

Machine Words in Isabelle/HOL

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Abstract

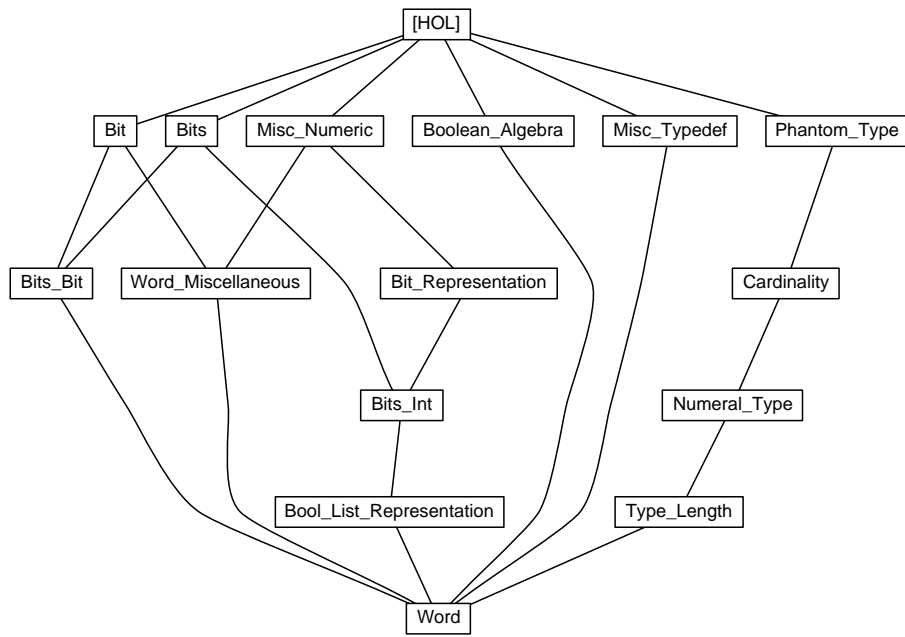
A formalisation of generic, fixed size machine words in Isabelle/HOL.
An earlier version of this formalisation is described in [1].

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1 A generic phantom type

```

theory Phantom-Type
imports Main
begin

datatype ('a, 'b) phantom = phantom (of-phantom: 'b)

lemma type-definition-phantom': type-definition of-phantom phantom UNIV
  ⟨proof⟩

lemma phantom-comp-of-phantom [simp]: phantom ∘ of-phantom = id
  and of-phantom-comp-phantom [simp]: of-phantom ∘ phantom = id
  ⟨proof⟩

syntax -Phantom :: type ⇒ logic ((1Phantom/(1'(-'))))
translations
  Phantom('t) => CONST phantom :: - ⇒ ('t, -) phantom

  ⟨ML⟩

lemma of-phantom-inject [simp]:
  of-phantom x = of-phantom y ⟷ x = y
  ⟨proof⟩

end

```

2 Cardinality of types

```

theory Cardinality
imports Phantom-Type
begin

```

2.1 Preliminary lemmas

```

lemma (in type-definition) univ:
  UNIV = Abs ' A
  ⟨proof⟩

lemma (in type-definition) card: card (UNIV :: 'b set) = card A
  ⟨proof⟩

lemma finite-range-Some: finite (range (Some :: 'a ⇒ 'a option)) = finite (UNIV
  :: 'a set)
  ⟨proof⟩

lemma infinite-literal: ¬ finite (UNIV :: String.literal set)
  ⟨proof⟩

```

2.2 Cardinalities of types

syntax *-type-card* :: *type* => *nat* ((1CARD/(1'(-))))

translations *CARD*('t) => *CONST card* (*CONST UNIV* :: 't set)

⟨ML⟩

lemma *card-prod* [*simp*]: $CARD('a \times 'b) = CARD('a) * CARD('b)$
 ⟨proof⟩

lemma *card-UNIV-sum*: $CARD('a + 'b) = (if\ CARD('a) \neq 0 \wedge CARD('b) \neq 0$
then $CARD('a) + CARD('b)$ *else* $0)$
 ⟨proof⟩

lemma *card-sum* [*simp*]: $CARD('a + 'b) = CARD('a::finite) + CARD('b::finite)$
 ⟨proof⟩

lemma *card-UNIV-option*: $CARD('a\ option) = (if\ CARD('a) = 0\ then\ 0\ else$
 $CARD('a) + 1)$
 ⟨proof⟩

lemma *card-option* [*simp*]: $CARD('a\ option) = Suc\ CARD('a::finite)$
 ⟨proof⟩

lemma *card-UNIV-set*: $CARD('a\ set) = (if\ CARD('a) = 0\ then\ 0\ else\ 2 \wedge CARD('a))$
 ⟨proof⟩

lemma *card-set* [*simp*]: $CARD('a\ set) = 2 \wedge CARD('a::finite)$
 ⟨proof⟩

lemma *card-nat* [*simp*]: $CARD(nat) = 0$
 ⟨proof⟩

lemma *card-fun*: $CARD('a \Rightarrow 'b) = (if\ CARD('a) \neq 0 \wedge CARD('b) \neq 0 \vee$
 $CARD('b) = 1\ then\ CARD('b) \wedge CARD('a)$ *else* $0)$
 ⟨proof⟩

corollary *finite-UNIV-fun*:

$finite\ (UNIV :: ('a \Rightarrow 'b)\ set) \longleftrightarrow$
 $finite\ (UNIV :: 'a\ set) \wedge finite\ (UNIV :: 'b\ set) \vee CARD('b) = 1$
 (is ?lhs \longleftrightarrow ?rhs)

⟨proof⟩

lemma *card-literal*: $CARD(String.literal) = 0$
 ⟨proof⟩

2.3 Classes with at least 1 and 2

Class *finite* already captures ”at least 1”

lemma *zero-less-card-finite* [*simp*]: $0 < \text{CARD}('a::\text{finite})$
 ⟨*proof*⟩

lemma *one-le-card-finite* [*simp*]: $\text{Suc } 0 \leq \text{CARD}('a::\text{finite})$
 ⟨*proof*⟩

Class for cardinality ”at least 2”

class *card2* = *finite* +
assumes *two-le-card*: $2 \leq \text{CARD}('a)$

lemma *one-less-card*: $\text{Suc } 0 < \text{CARD}('a::\text{card2})$
 ⟨*proof*⟩

lemma *one-less-int-card*: $1 < \text{int } \text{CARD}('a::\text{card2})$
 ⟨*proof*⟩

2.4 A type class for deciding finiteness of types

type-synonym *'a finite-UNIV* = (*'a*, *bool*) *phantom*

class *finite-UNIV* =
fixes *finite-UNIV* :: (*'a*, *bool*) *phantom*
assumes *finite-UNIV*: *finite-UNIV* = *Phantom*('a) (*finite* (*UNIV* :: 'a *set*))

lemma *finite-UNIV-code* [*code-unfold*]:
finite (*UNIV* :: 'a :: *finite-UNIV set*)
 \longleftrightarrow *of-phantom* (*finite-UNIV* :: 'a *finite-UNIV*)
 ⟨*proof*⟩

2.5 A type class for computing the cardinality of types

definition *is-list-UNIV* :: 'a *list* \Rightarrow *bool*
where *is-list-UNIV* *xs* = (let *c* = *CARD*('a) in if *c* = 0 then *False* else *size*
 (*remdups* *xs*) = *c*)

lemma *is-list-UNIV-iff*: *is-list-UNIV* *xs* \longleftrightarrow *set* *xs* = *UNIV*
 ⟨*proof*⟩

type-synonym 'a *card-UNIV* = (*'a*, *nat*) *phantom*

class *card-UNIV* = *finite-UNIV* +
fixes *card-UNIV* :: 'a *card-UNIV*
assumes *card-UNIV*: *card-UNIV* = *Phantom*('a) *CARD*('a)

2.6 Instantiations for *card-UNIV*

instantiation *nat* :: *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom*(*nat*) *False*
definition *card-UNIV* = *Phantom*(*nat*) 0
instance ⟨*proof*⟩

end

instantiation *int* :: *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom(int) False*
definition *card-UNIV* = *Phantom(int) 0*
instance *<proof>*
end

instantiation *natural* :: *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom(natural) False*
definition *card-UNIV* = *Phantom(natural) 0*
instance
<proof>
end

instantiation *integer* :: *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom(integer) False*
definition *card-UNIV* = *Phantom(integer) 0*
instance
<proof>
end

instantiation *list* :: (*type*) *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom('a list) False*
definition *card-UNIV* = *Phantom('a list) 0*
instance *<proof>*
end

instantiation *unit* :: *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom(unit) True*
definition *card-UNIV* = *Phantom(unit) 1*
instance *<proof>*
end

instantiation *bool* :: *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom(bool) True*
definition *card-UNIV* = *Phantom(bool) 2*
instance *<proof>*
end

instantiation *char* :: *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom(char) True*
definition *card-UNIV* = *Phantom(char) 256*
instance *<proof>*
end

instantiation *prod* :: (*finite-UNIV*, *finite-UNIV*) *finite-UNIV* **begin**
definition *finite-UNIV* = *Phantom('a × 'b)*
(of-phantom (finite-UNIV :: 'a finite-UNIV) ∧ of-phantom (finite-UNIV :: 'b


```

finite-UNIV)
instance <proof>
end

instantiation prod :: (card-UNIV, card-UNIV) card-UNIV begin
definition card-UNIV = Phantom('a × 'b)
  (of-phantom (card-UNIV :: 'a card-UNIV) * of-phantom (card-UNIV :: 'b card-UNIV))
instance <proof>
end

instantiation sum :: (finite-UNIV, finite-UNIV) finite-UNIV begin
definition finite-UNIV = Phantom('a + 'b)
  (of-phantom (finite-UNIV :: 'a finite-UNIV) ∧ of-phantom (finite-UNIV :: 'b
finite-UNIV))
instance
  <proof>
end

instantiation sum :: (card-UNIV, card-UNIV) card-UNIV begin
definition card-UNIV = Phantom('a + 'b)
  (let ca = of-phantom (card-UNIV :: 'a card-UNIV);
      cb = of-phantom (card-UNIV :: 'b card-UNIV)
      in if ca ≠ 0 ∧ cb ≠ 0 then ca + cb else 0)
instance <proof>
end

instantiation fun :: (finite-UNIV, card-UNIV) finite-UNIV begin
definition finite-UNIV = Phantom('a ⇒ 'b)
  (let cb = of-phantom (card-UNIV :: 'b card-UNIV)
      in cb = 1 ∨ of-phantom (finite-UNIV :: 'a finite-UNIV) ∧ cb ≠ 0)
instance
  <proof>
end

instantiation fun :: (card-UNIV, card-UNIV) card-UNIV begin
definition card-UNIV = Phantom('a ⇒ 'b)
  (let ca = of-phantom (card-UNIV :: 'a card-UNIV);
      cb = of-phantom (card-UNIV :: 'b card-UNIV)
      in if ca ≠ 0 ∧ cb ≠ 0 ∨ cb = 1 then cb ^ ca else 0)
instance <proof>
end

instantiation option :: (finite-UNIV) finite-UNIV begin
definition finite-UNIV = Phantom('a option) (of-phantom (finite-UNIV :: 'a
finite-UNIV))
instance <proof>
end

instantiation option :: (card-UNIV) card-UNIV begin

```

definition *card-UNIV* = *Phantom('a option)*
 (let *c* = *of-phantom (card-UNIV :: 'a card-UNIV)* in if *c* ≠ 0 then *Suc c* else 0)
instance ⟨*proof*⟩
end

instantiation *String.literal* :: *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom(String.literal) False*
definition *card-UNIV* = *Phantom(String.literal) 0*
instance
 ⟨*proof*⟩
end

instantiation *set* :: (*finite-UNIV*) *finite-UNIV* **begin**
definition *finite-UNIV* = *Phantom('a set) (of-phantom (finite-UNIV :: 'a finite-UNIV))*
instance ⟨*proof*⟩
end

instantiation *set* :: (*card-UNIV*) *card-UNIV* **begin**
definition *card-UNIV* = *Phantom('a set)*
 (let *c* = *of-phantom (card-UNIV :: 'a card-UNIV)* in if *c* = 0 then 0 else 2 ^ *c*)
instance ⟨*proof*⟩
end

lemma *UNIV-finite-1*: *UNIV* = *set [finite-1.a₁]*
 ⟨*proof*⟩

lemma *UNIV-finite-2*: *UNIV* = *set [finite-2.a₁, finite-2.a₂]*
 ⟨*proof*⟩

lemma *UNIV-finite-3*: *UNIV* = *set [finite-3.a₁, finite-3.a₂, finite-3.a₃]*
 ⟨*proof*⟩

lemma *UNIV-finite-4*: *UNIV* = *set [finite-4.a₁, finite-4.a₂, finite-4.a₃, finite-4.a₄]*
 ⟨*proof*⟩

lemma *UNIV-finite-5*:
UNIV = *set [finite-5.a₁, finite-5.a₂, finite-5.a₃, finite-5.a₄, finite-5.a₅]*
 ⟨*proof*⟩

instantiation *Enum.finite-1* :: *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom(Enum.finite-1) True*
definition *card-UNIV* = *Phantom(Enum.finite-1) 1*
instance
 ⟨*proof*⟩
end

instantiation *Enum.finite-2* :: *card-UNIV* **begin**
definition *finite-UNIV* = *Phantom(Enum.finite-2) True*
definition *card-UNIV* = *Phantom(Enum.finite-2) 2*

```

instance
  ⟨proof⟩
end

```

```

instantiation Enum.finite-3 :: card-UNIV begin
definition finite-UNIV = Phantom(Enum.finite-3) True
definition card-UNIV = Phantom(Enum.finite-3) 3
instance
  ⟨proof⟩
end

```

```

instantiation Enum.finite-4 :: card-UNIV begin
definition finite-UNIV = Phantom(Enum.finite-4) True
definition card-UNIV = Phantom(Enum.finite-4) 4
instance
  ⟨proof⟩
end

```

```

instantiation Enum.finite-5 :: card-UNIV begin
definition finite-UNIV = Phantom(Enum.finite-5) True
definition card-UNIV = Phantom(Enum.finite-5) 5
instance
  ⟨proof⟩
end

```

2.7 Code setup for sets

Implement $CARD('a)$ via $card-UNIV-class.card-UNIV$ and provide implementations for $finite$, $card$, $op \subseteq$, and $op =$ if the calling context already provides $finite-UNIV$ and $card-UNIV$ instances. If we implemented the latter always via $card-UNIV-class.card-UNIV$, we would require instances of essentially all element types, i.e., a lot of instantiation proofs and – at run time – possibly slow dictionary constructions.

```

context
begin

```

```

qualified definition card-UNIV' :: 'a card-UNIV
where [code del]: card-UNIV' = Phantom('a) CARD('a)

```

```

lemma CARD-code [code-unfold]:
  CARD('a) = of-phantom (card-UNIV' :: 'a card-UNIV)
  ⟨proof⟩

```

```

lemma card-UNIV'-code [code]:
  card-UNIV' = card-UNIV
  ⟨proof⟩

```

```

end

```

lemma *card-Compl*:

$finite\ A \implies card\ (-\ A) = card\ (UNIV :: 'a\ set) - card\ (A :: 'a\ set)$
 ⟨proof⟩

context **fixes** $xs :: 'a :: finite\ UNIV\ list$
begin

qualified definition $finite' :: 'a\ set \Rightarrow bool$
where [*simp*, *code del*, *code-abbrev*]: $finite' = finite$

lemma *finite'-code* [*code*]:
 $finite'\ (set\ xs) \longleftrightarrow True$
 $finite'\ (List.coset\ xs) \longleftrightarrow of\ phantom\ (finite\ UNIV :: 'a\ finite\ UNIV)$
 ⟨proof⟩

end

context **fixes** $xs :: 'a :: card\ UNIV\ list$
begin

qualified definition $card' :: 'a\ set \Rightarrow nat$
where [*simp*, *code del*, *code-abbrev*]: $card' = card$

lemma *card'-code* [*code*]:
 $card'\ (set\ xs) = length\ (remdups\ xs)$
 $card'\ (List.coset\ xs) = of\ phantom\ (card\ UNIV :: 'a\ card\ UNIV) - length\ (remdups\ xs)$
 ⟨proof⟩

definition $subset' :: 'a\ set \Rightarrow 'a\ set \Rightarrow bool$
where [*simp*, *code del*, *code-abbrev*]: $subset' = op\ \subseteq$

lemma *subset'-code* [*code*]:
 $subset'\ A\ (List.coset\ ys) \longleftrightarrow (\forall y \in set\ ys. y \notin A)$
 $subset'\ (set\ ys)\ B \longleftrightarrow (\forall y \in set\ ys. y \in B)$
 $subset'\ (List.coset\ xs)\ (set\ ys) \longleftrightarrow (let\ n = CARD('a)\ in\ n > 0 \wedge card\ (set\ (xs\ @\ ys)) = n)$
 ⟨proof⟩

definition $eq\ set :: 'a\ set \Rightarrow 'a\ set \Rightarrow bool$
where [*simp*, *code del*, *code-abbrev*]: $eq\ set = op\ =$

lemma *eq-set-code* [*code*]:

fixes ys
defines $rhs \equiv$
 $let\ n = CARD('a)$
 $in\ if\ n = 0\ then\ False\ else$
 $let\ xs' = remdups\ xs; ys' = remdups\ ys$
 $in\ length\ xs' + length\ ys' = n \wedge (\forall x \in set\ xs'. x \notin set\ ys') \wedge (\forall y \in set\ ys'.$
 $y \notin set\ xs')$
shows $eq\ set\ (List.coset\ xs)\ (set\ ys) \longleftrightarrow rhs$
and $eq\ set\ (set\ ys)\ (List.coset\ xs) \longleftrightarrow rhs$

```

and eq-set (set xs) (set ys)  $\longleftrightarrow$  ( $\forall x \in \text{set } xs. x \in \text{set } ys$ )  $\wedge$  ( $\forall y \in \text{set } ys. y \in \text{set } xs$ )
and eq-set (List.coset xs) (List.coset ys)  $\longleftrightarrow$  ( $\forall x \in \text{set } xs. x \in \text{set } ys$ )  $\wedge$  ( $\forall y \in \text{set } ys. y \in \text{set } xs$ )
<proof>

```

end

Provide more informative exceptions than Match for non-rewritten cases. If generated code raises one these exceptions, then a code equation calls the mentioned operator for an element type that is not an instance of *card-UNIV* and is therefore not implemented via *card-UNIV-class.card-UNIV*. Constrain the element type with sort *card-UNIV* to change this.

lemma *card-coset-error* [code]:

```

card (List.coset xs) =
  Code.abort (STR "card (List.coset -) requires type class instance card-UNIV")
    ( $\lambda$ -. card (List.coset xs))
<proof>

```

lemma *coset-subseteq-set-code* [code]:

```

List.coset xs  $\subseteq$  set ys  $\longleftrightarrow$ 
  (if xs = []  $\wedge$  ys = [] then False
   else Code.abort
     (STR "subset-eq (List.coset -) (List.set -) requires type class instance card-UNIV")
     ( $\lambda$ -. List.coset xs  $\subseteq$  set ys))
<proof>

```

notepad begin — test code setup

<proof>

end

end

3 Numeral Syntax for Types

theory *Numeral-Type*

imports *Cardinality*

begin

3.1 Numeral Types

typedef *num0* = *UNIV* :: *nat set* <proof>

typedef *num1* = *UNIV* :: *unit set* <proof>

typedef 'a *bit0* = {0 ..< 2 * int *CARD*('a::finite)}

<proof>

typedef 'a *bit1* = {0 ..< 1 + 2 * int *CARD*('a::finite)}

<proof>

lemma *card-num0* [*simp*]: $CARD (num0) = 0$
<proof>

lemma *infinite-num0*: $\neg finite (UNIV :: num0 set)$
<proof>

lemma *card-num1* [*simp*]: $CARD(num1) = 1$
<proof>

lemma *card-bit0* [*simp*]: $CARD('a bit0) = 2 * CARD('a::finite)$
<proof>

lemma *card-bit1* [*simp*]: $CARD('a bit1) = Suc (2 * CARD('a::finite))$
<proof>

instance *num1* :: *finite*
<proof>

instance *bit0* :: (*finite*) *card2*
<proof>

instance *bit1* :: (*finite*) *card2*
<proof>

3.2 Locales for modular arithmetic subtypes

locale *mod-type* =
fixes *n* :: *int*
and *Rep* :: 'a::{*zero,one,plus,times,uminus,minus*} \Rightarrow *int*
and *Abs* :: *int* \Rightarrow 'a::{*zero,one,plus,times,uminus,minus*}
assumes *type*: *type-definition Rep Abs {0..*n*}*
and *size1*: $1 < n$
and *zero-def*: $0 = Abs\ 0$
and *one-def*: $1 = Abs\ 1$
and *add-def*: $x + y = Abs ((Rep\ x + Rep\ y)\ mod\ n)$
and *mult-def*: $x * y = Abs ((Rep\ x * Rep\ y)\ mod\ n)$
and *diff-def*: $x - y = Abs ((Rep\ x - Rep\ y)\ mod\ n)$
and *minus-def*: $- x = Abs ((- Rep\ x)\ mod\ n)$
begin

lemma *size0*: $0 < n$
<proof>

lemmas *definitions* =
zero-def one-def add-def mult-def minus-def diff-def

lemma *Rep-less-n*: $Rep\ x < n$

<proof>

lemma *Rep-le-n*: $Rep\ x \leq n$

<proof>

lemma *Rep-inject-sym*: $x = y \longleftrightarrow Rep\ x = Rep\ y$

<proof>

lemma *Rep-inverse*: $Abs\ (Rep\ x) = x$

<proof>

lemma *Abs-inverse*: $m \in \{0..<n\} \implies Rep\ (Abs\ m) = m$

<proof>

lemma *Rep-Abs-mod*: $Rep\ (Abs\ (m\ mod\ n)) = m\ mod\ n$

<proof>

lemma *Rep-Abs-0*: $Rep\ (Abs\ 0) = 0$

<proof>

lemma *Rep-0*: $Rep\ 0 = 0$

<proof>

lemma *Rep-Abs-1*: $Rep\ (Abs\ 1) = 1$

<proof>

lemma *Rep-1*: $Rep\ 1 = 1$

<proof>

lemma *Rep-mod*: $Rep\ x\ mod\ n = Rep\ x$

<proof>

lemmas *Rep-simps* =

Rep-inject-sym Rep-inverse Rep-Abs-mod Rep-mod Rep-Abs-0 Rep-Abs-1

lemma *comm-ring-1*: *OFCLASS('a, comm-ring-1-class)*

<proof>

end

locale *mod-ring* = *mod-type n Rep Abs*

for *n* :: *int*

and *Rep* :: 'a::{*comm-ring-1*} \Rightarrow *int*

and *Abs* :: *int* \Rightarrow 'a::{*comm-ring-1*}

begin

lemma *of-nat-eq*: $of\ nat\ k = Abs\ (int\ k\ mod\ n)$

<proof>

lemma *of-int-eq*: $of\text{-}int\ z = Abs\ (z\ mod\ n)$
 ⟨*proof*⟩

lemma *Rep-numeral*:
 $Rep\ (numeral\ w) = numeral\ w\ mod\ n$
 ⟨*proof*⟩

lemma *iszero-numeral*:
 $iszero\ (numeral\ w::'a) \longleftrightarrow numeral\ w\ mod\ n = 0$
 ⟨*proof*⟩

lemma *cases*:
assumes 1: $\bigwedge z. \llbracket (x::'a) = of\text{-}int\ z; 0 \leq z; z < n \rrbracket \implies P$
shows P
 ⟨*proof*⟩

lemma *induct*:
 $(\bigwedge z. \llbracket 0 \leq z; z < n \rrbracket \implies P\ (of\text{-}int\ z)) \implies P\ (x::'a)$
 ⟨*proof*⟩

end

3.3 Ring class instances

Unfortunately *ring-1* instance is not possible for *num1*, since 0 and 1 are not distinct.

instantiation *num1* :: {*comm-ring,comm-monoid-mult,numeral*}
begin

lemma *num1-eq-iff*: $(x::num1) = (y::num1) \longleftrightarrow True$
 ⟨*proof*⟩

instance
 ⟨*proof*⟩

end

instantiation
bit0 and *bit1* :: (*finite*) {*zero,one,plus,times,uminus,minus*}
begin

definition *Abs-bit0'* :: $int \Rightarrow 'a\ bit0$ **where**
 $Abs\text{-}bit0'\ x = Abs\text{-}bit0\ (x\ mod\ int\ CARD('a\ bit0))$

definition *Abs-bit1'* :: $int \Rightarrow 'a\ bit1$ **where**
 $Abs\text{-}bit1'\ x = Abs\text{-}bit1\ (x\ mod\ int\ CARD('a\ bit1))$

definition $0 = Abs\text{-}bit0\ 0$

definition $1 = Abs\text{-}bit0\ 1$

definition $x + y = \text{Abs-bit0}' (\text{Rep-bit0 } x + \text{Rep-bit0 } y)$

definition $x * y = \text{Abs-bit0}' (\text{Rep-bit0 } x * \text{Rep-bit0 } y)$

definition $x - y = \text{Abs-bit0}' (\text{Rep-bit0 } x - \text{Rep-bit0 } y)$

definition $- x = \text{Abs-bit0}' (- \text{Rep-bit0 } x)$

definition $0 = \text{Abs-bit1 } 0$

definition $1 = \text{Abs-bit1 } 1$

definition $x + y = \text{Abs-bit1}' (\text{Rep-bit1 } x + \text{Rep-bit1 } y)$

definition $x * y = \text{Abs-bit1}' (\text{Rep-bit1 } x * \text{Rep-bit1 } y)$

definition $x - y = \text{Abs-bit1}' (\text{Rep-bit1 } x - \text{Rep-bit1 } y)$

definition $- x = \text{Abs-bit1}' (- \text{Rep-bit1 } x)$

instance $\langle \text{proof} \rangle$

end

interpretation *bit0*:

mod-type int CARD('a::finite bit0)

Rep-bit0 :: 'a::finite bit0 \Rightarrow int

Abs-bit0 :: int \Rightarrow 'a::finite bit0

$\langle \text{proof} \rangle$

interpretation *bit1*:

mod-type int CARD('a::finite bit1)

Rep-bit1 :: 'a::finite bit1 \Rightarrow int

Abs-bit1 :: int \Rightarrow 'a::finite bit1

$\langle \text{proof} \rangle$

instance *bit0* :: (*finite*) *comm-ring-1*

$\langle \text{proof} \rangle$

instance *bit1* :: (*finite*) *comm-ring-1*

$\langle \text{proof} \rangle$

interpretation *bit0*:

mod-ring int CARD('a::finite bit0)

Rep-bit0 :: 'a::finite bit0 \Rightarrow int

Abs-bit0 :: int \Rightarrow 'a::finite bit0

$\langle \text{proof} \rangle$

interpretation *bit1*:

mod-ring int CARD('a::finite bit1)

Rep-bit1 :: 'a::finite bit1 \Rightarrow int

Abs-bit1 :: int \Rightarrow 'a::finite bit1

$\langle \text{proof} \rangle$

Set up cases, induction, and arithmetic

lemmas *bit0-cases* [*case-names of-int*, *cases type: bit0*] = *bit0.cases*

lemmas *bit1-cases* [*case-names of-int*, *cases type: bit1*] = *bit1.cases*

lemmas *bit0-induct* [*case-names of-int, induct type: bit0*] = *bit0.induct*

lemmas *bit1-induct* [*case-names of-int, induct type: bit1*] = *bit1.induct*

lemmas *bit0-iszero-numeral* [*simp*] = *bit0.iszero-numeral*

lemmas *bit1-iszero-numeral* [*simp*] = *bit1.iszero-numeral*

lemmas [*simp*] = *eq-numeral-iff-iszero* [**where** *'a='a bit0*] **for** *dummy :: 'a::finite*

lemmas [*simp*] = *eq-numeral-iff-iszero* [**where** *'a='a bit1*] **for** *dummy :: 'a::finite*

3.4 Order instances

instantiation *bit0* and *bit1* :: (*finite*) *linorder* **begin**

definition $a < b \iff \text{Rep-bit0 } a < \text{Rep-bit0 } b$

definition $a \leq b \iff \text{Rep-bit0 } a \leq \text{Rep-bit0 } b$

definition $a < b \iff \text{Rep-bit1 } a < \text{Rep-bit1 } b$

definition $a \leq b \iff \text{Rep-bit1 } a \leq \text{Rep-bit1 } b$

instance

<proof>

end

lemma (**in preorder**) *tranclp-less: op <⁺⁺ = op <*

<proof>

instance *bit0* and *bit1* :: (*finite*) *wellorder*

<proof>

3.5 Code setup and type classes for code generation

Code setup for *num0* and *num1*

definition *Num0* :: *num0* **where** *Num0* = *Abs-num0 0*

code-datatype *Num0*

instantiation *num0* :: *equal* **begin**

definition *equal-num0* :: *num0* \Rightarrow *num0* \Rightarrow *bool*

where *equal-num0* = *op* =

instance *<proof>*

end

lemma *equal-num0-code* [*code*]:

equal-class.equal Num0 Num0 = *True*

<proof>

code-datatype *1* :: *num1*

instantiation *num1* :: *equal* **begin**

definition *equal-num1* :: *num1* \Rightarrow *num1* \Rightarrow *bool*

where *equal-num1* = *op* =

instance $\langle proof \rangle$
end

lemma *equal-num1-code* [code]:
equal-class.equal (1 :: num1) 1 = True
 $\langle proof \rangle$

instantiation num1 :: enum **begin**
definition *enum-class.enum* = [1 :: num1]
definition *enum-class.enum-all* P = P (1 :: num1)
definition *enum-class.enum-ex* P = P (1 :: num1)
instance
 $\langle proof \rangle$
end

instantiation num0 and num1 :: card-UNIV **begin**
definition *finite-UNIV* = Phantom(num0) False
definition *card-UNIV* = Phantom(num0) 0
definition *finite-UNIV* = Phantom(num1) True
definition *card-UNIV* = Phantom(num1) 1
instance
 $\langle proof \rangle$
end

Code setup for 'a bit0 and 'a bit1

declare
bit0.Rep-inverse[code abstype]
bit0.Rep-0[code abstract]
bit0.Rep-1[code abstract]

lemma *Abs-bit0'-code* [code abstract]:
Rep-bit0 (Abs-bit0' x :: 'a :: finite bit0) = x mod int (CARD('a bit0))
 $\langle proof \rangle$

lemma *inj-on-Abs-bit0*:
inj-on (Abs-bit0 :: int \Rightarrow 'a bit0) {0.. $2 * \text{int CARD}('a :: \text{finite})$ }
 $\langle proof \rangle$

declare
bit1.Rep-inverse[code abstype]
bit1.Rep-0[code abstract]
bit1.Rep-1[code abstract]

lemma *Abs-bit1'-code* [code abstract]:
Rep-bit1 (Abs-bit1' x :: 'a :: finite bit1) = x mod int (CARD('a bit1))
 $\langle proof \rangle$

lemma *inj-on-Abs-bit1*:
inj-on (Abs-bit1 :: int \Rightarrow 'a bit1) {0.. $1 + 2 * \text{int CARD}('a :: \text{finite})$ }

<proof>

instantiation *bit0* and *bit1* :: (*finite*) *equal* **begin**

definition *equal-class.equal* *x y* \longleftrightarrow *Rep-bit0* *x* = *Rep-bit0* *y*

definition *equal-class.equal* *x y* \longleftrightarrow *Rep-bit1* *x* = *Rep-bit1* *y*

instance

<proof>

end

instantiation *bit0* :: (*finite*) *enum* **begin**

definition (*enum-class.enum* :: '*a bit0 list*) = *map* (*Abs-bit0'* \circ *int*) (*upt 0* (*CARD*('a *bit0*)))

definition *enum-class.enum-all* *P* = (\forall *b* :: '*a bit0* \in *set enum-class.enum*. *P b*)

definition *enum-class.enum-ex* *P* = (\exists *b* :: '*a bit0* \in *set enum-class.enum*. *P b*)

instance

<proof>

end

instantiation *bit1* :: (*finite*) *enum* **begin**

definition (*enum-class.enum* :: '*a bit1 list*) = *map* (*Abs-bit1'* \circ *int*) (*upt 0* (*CARD*('a *bit1*)))

definition *enum-class.enum-all* *P* = (\forall *b* :: '*a bit1* \in *set enum-class.enum*. *P b*)

definition *enum-class.enum-ex* *P* = (\exists *b* :: '*a bit1* \in *set enum-class.enum*. *P b*)

instance

<proof>

end

instantiation *bit0* and *bit1* :: (*finite*) *finite-UNIV* **begin**

definition *finite-UNIV* = *Phantom*('a *bit0*) *True*

definition *finite-UNIV* = *Phantom*('a *bit1*) *True*

instance *<proof>*

end

instantiation *bit0* and *bit1* :: (*{finite,card-UNIV}*) *card-UNIV* **begin**

definition *card-UNIV* = *Phantom*('a *bit0*) (*2 * of-phantom* (*card-UNIV* :: '*a card-UNIV*))

definition *card-UNIV* = *Phantom*('a *bit1*) (*1 + 2 * of-phantom* (*card-UNIV* :: '*a card-UNIV*))

instance *<proof>*

end

3.6 Syntax

syntax

-*NumeralType* :: *num-token* => *type* (-)
 -*NumeralType0* :: *type* (0)
 -*NumeralType1* :: *type* (1)

translations

(*type*) 1 == (*type*) *num1*
 (*type*) 0 == (*type*) *num0*

⟨*ML*⟩

3.7 Examples

lemma *CARD*(0) = 0 ⟨*proof*⟩
lemma *CARD*(17) = 17 ⟨*proof*⟩
lemma $8 * 11 ^ 3 - 6 = (2::5)$ ⟨*proof*⟩

end

4 Assigning lengths to types by typeclasses

theory *Type-Length*

imports $\sim\sim$ /*src/HOL/Library/Numeral-Type*

begin

The aim of this is to allow any type as index type, but to provide a default instantiation for numeral types. This independence requires some duplication with the definitions in *Numeral-Type*.

class *len0* =
fixes *len-of* :: 'a *itself* \Rightarrow *nat*

Some theorems are only true on words with length greater 0.

class *len* = *len0* +
assumes *len-gt-0* [*iff*]: 0 < *len-of TYPE* ('a)

instantiation *num0* **and** *num1* :: *len0*

begin

definition

len-num0: *len-of* (*x*::*num0* *itself*) = 0

definition

len-num1: *len-of* (*x*::*num1* *itself*) = 1

instance ⟨*proof*⟩

end

instantiation *bit0* and *bit1* :: (*len0*) *len0*
begin

definition

len-bit0: *len-of* (*x*::'*a*::*len0* *bit0* *itself*) = 2 * *len-of* *TYPE* ('*a*)

definition

len-bit1: *len-of* (*x*::'*a*::*len0* *bit1* *itself*) = 2 * *len-of* *TYPE* ('*a*) + 1

instance ⟨*proof*⟩

end

lemmas *len-of-numeral-defs* [*simp*] = *len-num0* *len-num1* *len-bit0* *len-bit1*

instance *num1* :: *len* ⟨*proof*⟩

instance *bit0* :: (*len*) *len* ⟨*proof*⟩

instance *bit1* :: (*len0*) *len* ⟨*proof*⟩

end

5 Boolean Algebras

theory *Boolean-Algebra*

imports *Main*

begin

locale *boolean* =

fixes *conj* :: '*a* ⇒ '*a* ⇒ '*a* (**infixr** \sqcap 70)

fixes *disj* :: '*a* ⇒ '*a* ⇒ '*a* (**infixr** \sqcup 65)

fixes *compl* :: '*a* ⇒ '*a* (\sim - [81] 80)

fixes *zero* :: '*a* (**0**)

fixes *one* :: '*a* (**1**)

assumes *conj-assoc*: ($x \sqcap y$) $\sqcap z = x \sqcap (y \sqcap z)$

assumes *disj-assoc*: ($x \sqcup y$) $\sqcup z = x \sqcup (y \sqcup z)$

assumes *conj-commute*: $x \sqcap y = y \sqcap x$

assumes *disj-commute*: $x \sqcup y = y \sqcup x$

assumes *conj-disj-distrib*: $x \sqcap (y \sqcup z) = (x \sqcap y) \sqcup (x \sqcap z)$

assumes *disj-conj-distrib*: $x \sqcup (y \sqcap z) = (x \sqcup y) \sqcap (x \sqcup z)$

assumes *conj-one-right* [*simp*]: $x \sqcap \mathbf{1} = x$

assumes *disj-zero-right* [*simp*]: $x \sqcup \mathbf{0} = x$

assumes *conj-cancel-right* [*simp*]: $x \sqcap \sim x = \mathbf{0}$

assumes *disj-cancel-right* [*simp*]: $x \sqcup \sim x = \mathbf{1}$

begin

sublocale *conj*: *abel-semigroup* *conj*

⟨*proof*⟩

sublocale *disj*: *abel-semigroup disj*
 ⟨*proof*⟩

lemmas *conj-left-commute* = *conj.left-commute*

lemmas *disj-left-commute* = *disj.left-commute*

lemmas *conj-ac* = *conj.assoc conj.commute conj.left-commute*

lemmas *disj-ac* = *disj.assoc disj.commute disj.left-commute*

lemma *dual*: *boolean disj compl one zero*
 ⟨*proof*⟩

5.1 Complement

lemma *complement-unique*:

assumes 1: $a \sqcap x = \mathbf{0}$

assumes 2: $a \sqcup x = \mathbf{1}$

assumes 3: $a \sqcap y = \mathbf{0}$

assumes 4: $a \sqcup y = \mathbf{1}$

shows $x = y$

⟨*proof*⟩

lemma *compl-unique*: $\llbracket x \sqcap y = \mathbf{0}; x \sqcup y = \mathbf{1} \rrbracket \implies \sim x = y$
 ⟨*proof*⟩

lemma *double-compl [simp]*: $\sim(\sim x) = x$
 ⟨*proof*⟩

lemma *compl-eq-compl-iff [simp]*: $(\sim x = \sim y) = (x = y)$
 ⟨*proof*⟩

5.2 Conjunction

lemma *conj-absorb [simp]*: $x \sqcap x = x$
 ⟨*proof*⟩

lemma *conj-zero-right [simp]*: $x \sqcap \mathbf{0} = \mathbf{0}$
 ⟨*proof*⟩

lemma *compl-one [simp]*: $\sim \mathbf{1} = \mathbf{0}$
 ⟨*proof*⟩

lemma *conj-zero-left [simp]*: $\mathbf{0} \sqcap x = \mathbf{0}$
 ⟨*proof*⟩

lemma *conj-one-left [simp]*: $\mathbf{1} \sqcap x = x$
 ⟨*proof*⟩

lemma *conj-cancel-left [simp]*: $\sim x \sqcap x = \mathbf{0}$

<proof>

lemma *conj-left-absorb* [*simp*]: $x \sqcap (x \sqcap y) = x \sqcap y$
<proof>

lemma *conj-disj-distrib2*:
 $(y \sqcup z) \sqcap x = (y \sqcap x) \sqcup (z \sqcap x)$
<proof>

lemmas *conj-disj-distrib* =
conj-disj-distrib conj-disj-distrib2

5.3 Disjunction

lemma *disj-absorb* [*simp*]: $x \sqcup x = x$
<proof>

lemma *disj-one-right* [*simp*]: $x \sqcup \mathbf{1} = \mathbf{1}$
<proof>

lemma *compl-zero* [*simp*]: $\sim \mathbf{0} = \mathbf{1}$
<proof>

lemma *disj-zero-left* [*simp*]: $\mathbf{0} \sqcup x = x$
<proof>

lemma *disj-one-left* [*simp*]: $\mathbf{1} \sqcup x = \mathbf{1}$
<proof>

lemma *disj-cancel-left* [*simp*]: $\sim x \sqcup x = \mathbf{1}$
<proof>

lemma *disj-left-absorb* [*simp*]: $x \sqcup (x \sqcup y) = x \sqcup y$
<proof>

lemma *disj-conj-distrib2*:
 $(y \sqcap z) \sqcup x = (y \sqcup x) \sqcap (z \sqcup x)$
<proof>

lemmas *disj-conj-distrib* =
disj-conj-distrib disj-conj-distrib2

5.4 De Morgan’s Laws

lemma *de-Morgan-conj* [*simp*]: $\sim (x \sqcap y) = \sim x \sqcup \sim y$
<proof>

lemma *de-Morgan-disj* [*simp*]: $\sim (x \sqcup y) = \sim x \sqcap \sim y$
<proof>

end

5.5 Symmetric Difference

locale *boolean-xor* = *boolean* +
fixes *xor* :: 'a ⇒ 'a ⇒ 'a (**infixr** ⊕ 65)
assumes *xor-def*: $x \oplus y = (x \sqcap \sim y) \sqcup (\sim x \sqcap y)$
begin

sublocale *xor*: *abel-semigroup xor*
 ⟨*proof*⟩

lemmas *xor-assoc* = *xor.assoc*
lemmas *xor-commute* = *xor.commute*
lemmas *xor-left-commute* = *xor.left-commute*

lemmas *xor-ac* = *xor.assoc xor.commute xor.left-commute*

lemma *xor-def2*:
 $x \oplus y = (x \sqcup y) \sqcap (\sim x \sqcup \sim y)$
 ⟨*proof*⟩

lemma *xor-zero-right* [*simp*]: $x \oplus \mathbf{0} = x$
 ⟨*proof*⟩

lemma *xor-zero-left* [*simp*]: $\mathbf{0} \oplus x = x$
 ⟨*proof*⟩

lemma *xor-one-right* [*simp*]: $x \oplus \mathbf{1} = \sim x$
 ⟨*proof*⟩

lemma *xor-one-left* [*simp*]: $\mathbf{1} \oplus x = \sim x$
 ⟨*proof*⟩

lemma *xor-self* [*simp*]: $x \oplus x = \mathbf{0}$
 ⟨*proof*⟩

lemma *xor-left-self* [*simp*]: $x \oplus (x \oplus y) = y$
 ⟨*proof*⟩

lemma *xor-compl-left* [*simp*]: $\sim x \oplus y = \sim (x \oplus y)$
 ⟨*proof*⟩

lemma *xor-compl-right* [*simp*]: $x \oplus \sim y = \sim (x \oplus y)$
 ⟨*proof*⟩

lemma *xor-cancel-right*: $x \oplus \sim x = \mathbf{1}$
 ⟨*proof*⟩

lemma *xor-cancel-left*: $\sim x \oplus x = \mathbf{1}$
 ⟨*proof*⟩

lemma *conj-xor-distrib*: $x \sqcap (y \oplus z) = (x \sqcap y) \oplus (x \sqcap z)$
 ⟨*proof*⟩

lemma *conj-xor-distrib2*: $(y \oplus z) \sqcap x = (y \sqcap x) \oplus (z \sqcap x)$
 ⟨*proof*⟩

lemmas *conj-xor-distrib* = *conj-xor-distrib conj-xor-distrib2*

end

end

6 Syntactic classes for bitwise operations

theory *Bits*
imports *Main*
begin

class *bit* =
fixes *bitNOT* :: 'a ⇒ 'a (*NOT* - [70] 71)
and *bitAND* :: 'a ⇒ 'a ⇒ 'a (**infixr** *AND* 64)
and *bitOR* :: 'a ⇒ 'a ⇒ 'a (**infixr** *OR* 59)
and *bitXOR* :: 'a ⇒ 'a ⇒ 'a (**infixr** *XOR* 59)

We want the bitwise operations to bind slightly weaker than + and −, but ~ to bind slightly stronger than *.

Testing and shifting operations.

class *bits* = *bit* +
fixes *test-bit* :: 'a ⇒ nat ⇒ bool (**infixl** !! 100)
and *lsb* :: 'a ⇒ bool
and *set-bit* :: 'a ⇒ nat ⇒ bool ⇒ 'a
and *set-bits* :: (nat ⇒ bool) ⇒ 'a (**binder** *BITS* 10)
and *shiffl* :: 'a ⇒ nat ⇒ 'a (**infixl** << 55)
and *shiftr* :: 'a ⇒ nat ⇒ 'a (**infixl** >> 55)

class *bitss* = *bits* +
fixes *msb* :: 'a ⇒ bool

end

7 The Field of Integers mod 2

theory *Bit*
imports *Main*

begin

7.1 Bits as a datatype

typedef *bit* = UNIV :: bool set
morphisms set *Bit*
 ⟨*proof*⟩

instantiation *bit* :: {*zero*, *one*}
begin

definition *zero-bit-def*:
 0 = Bit False

definition *one-bit-def*:
 1 = Bit True

instance ⟨*proof*⟩

end

old-rep-datatype 0::bit 1::bit
 ⟨*proof*⟩

lemma *Bit-set-eq* [*simp*]:
 Bit (set b) = b
 ⟨*proof*⟩

lemma *set-Bit-eq* [*simp*]:
 set (Bit P) = P
 ⟨*proof*⟩

lemma *bit-eq-iff*:
 $x = y \longleftrightarrow (\text{set } x \longleftrightarrow \text{set } y)$
 ⟨*proof*⟩

lemma *Bit-inject* [*simp*]:
 $\text{Bit } P = \text{Bit } Q \longleftrightarrow (P \longleftrightarrow Q)$
 ⟨*proof*⟩

lemma *set* [*iff*]:
 $\neg \text{set } 0$
 $\text{set } 1$
 ⟨*proof*⟩

lemma [*code*]:
 $\text{set } 0 \longleftrightarrow \text{False}$
 $\text{set } 1 \longleftrightarrow \text{True}$
 ⟨*proof*⟩

lemma *set-iff*:
 $set\ b \longleftrightarrow b = 1$
 ⟨proof⟩

lemma *bit-eq-iff-set*:
 $b = 0 \longleftrightarrow \neg\ set\ b$
 $b = 1 \longleftrightarrow set\ b$
 ⟨proof⟩

lemma *Bit* [*simp*, *code*]:
 $Bit\ False = 0$
 $Bit\ True = 1$
 ⟨proof⟩

lemma *bit-not-0-iff* [*iff*]:
 $(x::bit) \neq 0 \longleftrightarrow x = 1$
 ⟨proof⟩

lemma *bit-not-1-iff* [*iff*]:
 $(x::bit) \neq 1 \longleftrightarrow x = 0$
 ⟨proof⟩

lemma [*code*]:
 $HOL.equal\ 0\ b \longleftrightarrow \neg\ set\ b$
 $HOL.equal\ 1\ b \longleftrightarrow set\ b$
 ⟨proof⟩

7.2 Type *bit* forms a field

instantiation *bit* :: *field*
begin

definition *plus-bit-def*:
 $x + y = case-bit\ y\ (case-bit\ 1\ 0\ y)\ x$

definition *times-bit-def*:
 $x * y = case-bit\ 0\ y\ x$

definition *uminus-bit-def* [*simp*]:
 $- x = (x :: bit)$

definition *minus-bit-def* [*simp*]:
 $x - y = (x + y :: bit)$

definition *inverse-bit-def* [*simp*]:
 $inverse\ x = (x :: bit)$

definition *divide-bit-def* [*simp*]:

$$x \text{ div } y = (x * y :: \text{bit})$$

lemmas *field-bit-defs* =
plus-bit-def times-bit-def minus-bit-def uminus-bit-def
divide-bit-def inverse-bit-def

instance

<proof>

end

lemma *bit-add-self*: $x + x = (0 :: \text{bit})$

<proof>

lemma *bit-mult-eq-1-iff* [*simp*]: $x * y = (1 :: \text{bit}) \longleftrightarrow x = 1 \wedge y = 1$

<proof>

Not sure whether the next two should be simp rules.

lemma *bit-add-eq-0-iff*: $x + y = (0 :: \text{bit}) \longleftrightarrow x = y$

<proof>

lemma *bit-add-eq-1-iff*: $x + y = (1 :: \text{bit}) \longleftrightarrow x \neq y$

<proof>

7.3 Numerals at type *bit*

All numerals reduce to either 0 or 1.

lemma *bit-minus1* [*simp*]: $- 1 = (1 :: \text{bit})$

<proof>

lemma *bit-neg-numeral* [*simp*]: $(- \text{ numeral } w :: \text{bit}) = \text{ numeral } w$

<proof>

lemma *bit-numeral-even* [*simp*]: $\text{ numeral } (\text{Num.Bit0 } w) = (0 :: \text{bit})$

<proof>

lemma *bit-numeral-odd* [*simp*]: $\text{ numeral } (\text{Num.Bit1 } w) = (1 :: \text{bit})$

<proof>

7.4 Conversion from *bit*

context *zero-neq-one*

begin

definition *of-bit* :: $\text{bit} \Rightarrow 'a$

where

of-bit b = case-bit 0 1 b

lemma *of-bit-eq* [*simp*, *code*]:

of-bit 0 = 0
of-bit 1 = 1
 ⟨*proof*⟩

lemma *of-bit-eq-iff*:
of-bit x = of-bit y \leftrightarrow *x = y*
 ⟨*proof*⟩

end

context *semiring-1*
begin

lemma *of-nat-of-bit-eq*:
of-nat (of-bit b) = of-bit b
 ⟨*proof*⟩

end

context *ring-1*
begin

lemma *of-int-of-bit-eq*:
of-int (of-bit b) = of-bit b
 ⟨*proof*⟩

end

hide-const (open) *set*

end

8 Bit operations in \mathcal{Z}_ϵ

theory *Bits-Bit*
imports *Bits* $\sim\sim$ /src/HOL/Library/Bit
begin

instantiation *bit* :: *bit*
begin

primrec *bitNOT-bit* **where**
NOT 0 = (1::bit)
 | *NOT 1 = (0::bit)*

primrec *bitAND-bit* **where**
0 AND y = (0::bit)
 | *1 AND y = (y::bit)*

primrec *bitOR-bit* **where**

$0 \text{ OR } y = (y::\text{bit})$
 $| 1 \text{ OR } y = (1::\text{bit})$

primrec *bitXOR-bit* **where**

$0 \text{ XOR } y = (y::\text{bit})$
 $| 1 \text{ XOR } y = (\text{NOT } y :: \text{bit})$

instance $\langle \text{proof} \rangle$

end

lemmas *bit-simps* =

bitNOT-bit.simps bitAND-bit.simps bitOR-bit.simps bitXOR-bit.simps

lemma *bit-extra-simps* [*simp*]:

$x \text{ AND } 0 = (0::\text{bit})$
 $x \text{ AND } 1 = (x::\text{bit})$
 $x \text{ OR } 1 = (1::\text{bit})$
 $x \text{ OR } 0 = (x::\text{bit})$
 $x \text{ XOR } 1 = \text{NOT } (x::\text{bit})$
 $x \text{ XOR } 0 = (x::\text{bit})$
 $\langle \text{proof} \rangle$

lemma *bit-ops-comm*:

$(x::\text{bit}) \text{ AND } y = y \text{ AND } x$
 $(x::\text{bit}) \text{ OR } y = y \text{ OR } x$
 $(x::\text{bit}) \text{ XOR } y = y \text{ XOR } x$
 $\langle \text{proof} \rangle$

lemma *bit-ops-same* [*simp*]:

$(x::\text{bit}) \text{ AND } x = x$
 $(x::\text{bit}) \text{ OR } x = x$
 $(x::\text{bit}) \text{ XOR } x = 0$
 $\langle \text{proof} \rangle$

lemma *bit-not-not* [*simp*]: $\text{NOT } (\text{NOT } (x::\text{bit})) = x$

$\langle \text{proof} \rangle$

lemma *bit-or-def*: $(b::\text{bit}) \text{ OR } c = \text{NOT } (\text{NOT } b \text{ AND } \text{NOT } c)$

$\langle \text{proof} \rangle$

lemma *bit-xor-def*: $(b::\text{bit}) \text{ XOR } c = (b \text{ AND } \text{NOT } c) \text{ OR } (\text{NOT } b \text{ AND } c)$

$\langle \text{proof} \rangle$

lemma *bit-NOT-eq-1-iff* [*simp*]: $\text{NOT } (b::\text{bit}) = 1 \iff b = 0$

$\langle \text{proof} \rangle$

lemma *bit-AND-eq-1-iff* [*simp*]: $(a::\text{bit}) \text{ AND } b = 1 \iff a = 1 \wedge b = 1$

<proof>

end

9 Useful Numerical Lemmas

theory *Misc-Numeric*

imports *Main*

begin

lemma *mod-2-neq-1-eq-eq-0*:

fixes $k :: int$

shows $k \bmod 2 \neq 1 \longleftrightarrow k \bmod 2 = 0$

<proof>

lemma *z1pmod2*:

fixes $b :: int$

shows $(2 * b + 1) \bmod 2 = (1 :: int)$

<proof>

lemma *diff-le-eq'*:

$a - b \leq c \longleftrightarrow a \leq b + (c :: int)$

<proof>

lemma *emep1*:

fixes $n d :: int$

shows $even\ n \implies even\ d \implies 0 \leq d \implies (n + 1) \bmod d = (n \bmod d) + 1$

<proof>

lemma *int-mod-ge*:

$a < n \implies 0 < (n :: int) \implies a \leq a \bmod n$

<proof>

lemma *int-mod-ge'*:

$b < 0 \implies 0 < (n :: int) \implies b + n \leq b \bmod n$

<proof>

lemma *int-mod-le'*:

$(0 :: int) \leq b - n \implies b \bmod n \leq b - n$

<proof>

lemma *zless2*:

$0 < (2 :: int)$

<proof>

lemma *zless2p*:

$0 < (2 ^ n :: int)$

<proof>


```

lemma zle2p:
   $0 \leq (2 \wedge n :: int)$ 
  <proof>

lemma m1mod2k:
   $-1 \bmod 2 \wedge n = (2 \wedge n - 1 :: int)$ 
  <proof>

lemma p1mod22k':
  fixes  $b :: int$ 
  shows  $(1 + 2 * b) \bmod (2 * 2 \wedge n) = 1 + 2 * (b \bmod 2 \wedge n)$ 
  <proof>

lemma p1mod22k:
  fixes  $b :: int$ 
  shows  $(2 * b + 1) \bmod (2 * 2 \wedge n) = 2 * (b \bmod 2 \wedge n) + 1$ 
  <proof>

lemma int-mod-lem:
   $(0 :: int) < n ==> (0 \leq b \ \& \ b < n) = (b \bmod n = b)$ 
  <proof>

end

```

10 Integers as implicit bit strings

```

theory Bit-Representation
imports Misc-Numeric
begin

```

10.1 Constructors and destructors for binary integers

```

definition Bit ::  $int \Rightarrow bool \Rightarrow int$  (infixl BIT 90)

```

```

where

```

```

   $k \text{ BIT } b = (\text{if } b \text{ then } 1 \text{ else } 0) + k + k$ 

```

```

lemma Bit-B0:

```

```

   $k \text{ BIT } \text{False} = k + k$ 

```

```

  <proof>

```

```

lemma Bit-B1:

```

```

   $k \text{ BIT } \text{True} = k + k + 1$ 

```

```

  <proof>

```

```

lemma Bit-B0-2t:  $k \text{ BIT } \text{False} = 2 * k$ 

```

```

  <proof>

```

```

lemma Bit-B1-2t:  $k \text{ BIT } \text{True} = 2 * k + 1$ 

```

```

  <proof>

```

definition *bin-last* :: *int* \Rightarrow *bool*

where

$$\text{bin-last } w \longleftrightarrow w \bmod 2 = 1$$

lemma *bin-last-odd*:

$$\text{bin-last} = \text{odd}$$

<proof>

definition *bin-rest* :: *int* \Rightarrow *int*

where

$$\text{bin-rest } w = w \text{ div } 2$$

lemma *bin-rl-simp* [*simp*]:

$$\text{bin-rest } w \text{ BIT } \text{bin-last } w = w$$

<proof>

lemma *bin-rest-BIT* [*simp*]: *bin-rest* (*x BIT b*) = *x*

<proof>

lemma *bin-last-BIT* [*simp*]: *bin-last* (*x BIT b*) = *b*

<proof>

lemma *BIT-eq-iff* [*iff*]: *u BIT b = v BIT c* \longleftrightarrow *u = v* \wedge *b = c*

<proof>

lemma *BIT-bin-simps* [*simp*]:

$$\text{numeral } k \text{ BIT } \text{False} = \text{numeral } (\text{Num.Bit0 } k)$$

$$\text{numeral } k \text{ BIT } \text{True} = \text{numeral } (\text{Num.Bit1 } k)$$

$$(- \text{ numeral } k) \text{ BIT } \text{False} = - \text{ numeral } (\text{Num.Bit0 } k)$$

$$(- \text{ numeral } k) \text{ BIT } \text{True} = - \text{ numeral } (\text{Num.BitM } k)$$

<proof>

lemma *BIT-special-simps* [*simp*]:

$$\text{shows } 0 \text{ BIT } \text{False} = 0 \text{ and } 0 \text{ BIT } \text{True} = 1$$

$$\text{and } 1 \text{ BIT } \text{False} = 2 \text{ and } 1 \text{ BIT } \text{True} = 3$$

$$\text{and } (- 1) \text{ BIT } \text{False} = - 2 \text{ and } (- 1) \text{ BIT } \text{True} = - 1$$

<proof>

lemma *Bit-eq-0-iff*: *w BIT b = 0* \longleftrightarrow *w = 0* \wedge $\neg b$

<proof>

lemma *Bit-eq-m1-iff*: *w BIT b = -1* \longleftrightarrow *w = -1* \wedge *b*

<proof>

lemma *BitM-inc*: *Num.BitM* (*Num.inc w*) = *Num.Bit1 w*

<proof>

lemma *expand-BIT*:

$\text{numeral } (\text{Num.Bit0 } w) = \text{numeral } w \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } w) = \text{numeral } w \text{ BIT True}$
 $\neg \text{numeral } (\text{Num.Bit0 } w) = (\neg \text{numeral } w) \text{ BIT False}$
 $\neg \text{numeral } (\text{Num.Bit1 } w) = (\neg \text{numeral } (w + \text{Num.One})) \text{ BIT True}$
 ⟨proof⟩

lemma *bin-last-numeral-simps* [simp]:

$\neg \text{bin-last } 0$
 $\text{bin-last } 1$
 $\text{bin-last } (- 1)$
 bin-last Numeral1
 $\neg \text{bin-last } (\text{numeral } (\text{Num.Bit0 } w))$
 $\text{bin-last } (\text{numeral } (\text{Num.Bit1 } w))$
 $\neg \text{bin-last } (\neg \text{numeral } (\text{Num.Bit0 } w))$
 $\text{bin-last } (\neg \text{numeral } (\text{Num.Bit1 } w))$
 ⟨proof⟩

lemma *bin-rest-numeral-simps* [simp]:

$\text{bin-rest } 0 = 0$
 $\text{bin-rest } 1 = 0$
 $\text{bin-rest } (- 1) = - 1$
 $\text{bin-rest Numeral1} = 0$
 $\text{bin-rest } (\text{numeral } (\text{Num.Bit0 } w)) = \text{numeral } w$
 $\text{bin-rest } (\text{numeral } (\text{Num.Bit1 } w)) = \text{numeral } w$
 $\text{bin-rest } (\neg \text{numeral } (\text{Num.Bit0 } w)) = \neg \text{numeral } w$
 $\text{bin-rest } (\neg \text{numeral } (\text{Num.Bit1 } w)) = \neg \text{numeral } (w + \text{Num.One})$
 ⟨proof⟩

lemma *less-Bits*:

$v \text{ BIT } b < w \text{ BIT } c \longleftrightarrow v < w \vee v \leq w \wedge \neg b \wedge c$
 ⟨proof⟩

lemma *le-Bits*:

$v \text{ BIT } b \leq w \text{ BIT } c \longleftrightarrow v < w \vee v \leq w \wedge (\neg b \vee c)$
 ⟨proof⟩

lemma *pred-BIT-simps* [simp]:

$x \text{ BIT False } - 1 = (x - 1) \text{ BIT True}$
 $x \text{ BIT True } - 1 = x \text{ BIT False}$
 ⟨proof⟩

lemma *succ-BIT-simps* [simp]:

$x \text{ BIT False } + 1 = x \text{ BIT True}$
 $x \text{ BIT True } + 1 = (x + 1) \text{ BIT False}$
 ⟨proof⟩

lemma *add-BIT-simps* [simp]:

$x \text{ BIT False } + y \text{ BIT False} = (x + y) \text{ BIT False}$
 $x \text{ BIT False } + y \text{ BIT True} = (x + y) \text{ BIT True}$

$x \text{ BIT True} + y \text{ BIT False} = (x + y) \text{ BIT True}$
 $x \text{ BIT True} + y \text{ BIT True} = (x + y + 1) \text{ BIT False}$
 ⟨proof⟩

lemma *mult-BIT-simps* [simp]:
 $x \text{ BIT False} * y = (x * y) \text{ BIT False}$
 $x * y \text{ BIT False} = (x * y) \text{ BIT False}$
 $x \text{ BIT True} * y = (x * y) \text{ BIT False} + y$
 ⟨proof⟩

lemma *B-mod-2'*:
 $X = 2 \implies (w \text{ BIT True}) \text{ mod } X = 1 \ \& \ (w \text{ BIT False}) \text{ mod } X = 0$
 ⟨proof⟩

lemma *bin-ex-rl*: $EX \ w \ b. \ w \text{ BIT } b = \text{bin}$
 ⟨proof⟩

lemma *bin-exhaust*:
assumes $Q: \bigwedge x \ b. \ \text{bin} = x \text{ BIT } b \implies Q$
shows Q
 ⟨proof⟩

primrec *bin-nth* **where**
 $Z: \text{bin-nth } w \ 0 \longleftrightarrow \text{bin-last } w$
 $| \text{Suc}: \text{bin-nth } w \ (\text{Suc } n) \longleftrightarrow \text{bin-nth } (\text{bin-rest } w) \ n$

lemma *bin-abs-lem*:
 $\text{bin} = (w \text{ BIT } b) \implies \text{bin} \sim = -1 \dashrightarrow \text{bin} \sim = 0 \dashrightarrow$
 $\text{nat } |w| < \text{nat } |\text{bin}|$
 ⟨proof⟩

lemma *bin-induct*:
assumes *PPls*: $P \ 0$
and *PMin*: $P \ (- \ 1)$
and *PBit*: $!! \text{bin } \text{bit}. \ P \ \text{bin} \implies P \ (\text{bin } \text{BIT } \text{bit})$
shows $P \ \text{bin}$
 ⟨proof⟩

lemma *Bit-div2* [simp]: $(w \text{ BIT } b) \text{ div } 2 = w$
 ⟨proof⟩

lemma *bin-nth-eq-iff*:
 $\text{bin-nth } x = \text{bin-nth } y \longleftrightarrow x = y$
 ⟨proof⟩

lemmas *bin-eqI* = *ext* [THEN *bin-nth-eq-iff*] [THEN *iffD1*]

lemma *bin-eq-iff*:
 $x = y \longleftrightarrow (\forall n. \ \text{bin-nth } x \ n = \text{bin-nth } y \ n)$

<proof>

lemma *bin-nth-zero* [*simp*]: $\neg \text{bin-nth } 0 \ n$
<proof>

lemma *bin-nth-1* [*simp*]: $\text{bin-nth } 1 \ n \longleftrightarrow n = 0$
<proof>

lemma *bin-nth-minus1* [*simp*]: $\text{bin-nth } (-1) \ n$
<proof>

lemma *bin-nth-0-BIT*: $\text{bin-nth } (w \ \text{BIT } b) \ 0 \longleftrightarrow b$
<proof>

lemma *bin-nth-Suc-BIT*: $\text{bin-nth } (w \ \text{BIT } b) \ (\text{Suc } n) = \text{bin-nth } w \ n$
<proof>

lemma *bin-nth-minus* [*simp*]: $0 < n \implies \text{bin-nth } (w \ \text{BIT } b) \ n = \text{bin-nth } w \ (n - 1)$
<proof>

lemma *bin-nth-numeral*:
 $\text{bin-rest } x = y \implies \text{bin-nth } x \ (\text{numeral } n) = \text{bin-nth } y \ (\text{pred-numeral } n)$
<proof>

lemmas *bin-nth-numeral-simps* [*simp*] =
 $\text{bin-nth-numeral } [\text{OF } \text{bin-rest-numeral-simps } (2)]$
 $\text{bin-nth-numeral } [\text{OF } \text{bin-rest-numeral-simps } (5)]$
 $\text{bin-nth-numeral } [\text{OF } \text{bin-rest-numeral-simps } (6)]$
 $\text{bin-nth-numeral } [\text{OF } \text{bin-rest-numeral-simps } (7)]$
 $\text{bin-nth-numeral } [\text{OF } \text{bin-rest-numeral-simps } (8)]$

lemmas *bin-nth-simps* =
 $\text{bin-nth.Z } \text{bin-nth.Suc } \text{bin-nth-zero } \text{bin-nth-minus1}$
 $\text{bin-nth-numeral-simps}$

10.2 Truncating binary integers

definition *bin-sign* :: $\text{int} \Rightarrow \text{int}$

where

bin-sign-def: $\text{bin-sign } k = (\text{if } k \geq 0 \text{ then } 0 \text{ else } -1)$

lemma *bin-sign-simps* [*simp*]:
 $\text{bin-sign } 0 = 0$
 $\text{bin-sign } 1 = 0$
 $\text{bin-sign } (-1) = -1$
 $\text{bin-sign } (\text{numeral } k) = 0$
 $\text{bin-sign } (- \text{numeral } k) = -1$
 $\text{bin-sign } (w \ \text{BIT } b) = \text{bin-sign } w$

<proof>

lemma *bin-sign-rest* [simp]:
 $\text{bin-sign } (\text{bin-rest } w) = \text{bin-sign } w$
<proof>

primrec *bintrunc* :: $\text{nat} \Rightarrow \text{int} \Rightarrow \text{int}$ **where**
 $Z : \text{bintrunc } 0 \text{ bin} = 0$
 $| \text{Suc } n : \text{bintrunc } (\text{Suc } n) \text{ bin} = \text{bintrunc } n (\text{bin-rest } \text{bin}) \text{ BIT } (\text{bin-last } \text{bin})$

primrec *sbintrunc* :: $\text{nat} \Rightarrow \text{int} \Rightarrow \text{int}$ **where**
 $Z : \text{sbintrunc } 0 \text{ bin} = (\text{if } \text{bin-last } \text{bin} \text{ then } -1 \text{ else } 0)$
 $| \text{Suc } n : \text{sbintrunc } (\text{Suc } n) \text{ bin} = \text{sbintrunc } n (\text{bin-rest } \text{bin}) \text{ BIT } (\text{bin-last } \text{bin})$

lemma *sign-bintr*: $\text{bin-sign } (\text{bintrunc } n \ w) = 0$
<proof>

lemma *bintrunc-mod2p*: $\text{bintrunc } n \ w = (w \bmod 2^n)$
<proof>

lemma *sbintrunc-mod2p*: $\text{sbintrunc } n \ w = (w + 2^n) \bmod 2^{(\text{Suc } n)} - 2^n$
<proof>

10.3 Simplifications for (s)bintrunc

lemma *bintrunc-n-0* [simp]: $\text{bintrunc } n \ 0 = 0$
<proof>

lemma *sbintrunc-n-0* [simp]: $\text{sbintrunc } n \ 0 = 0$
<proof>

lemma *sbintrunc-n-minus1* [simp]: $\text{sbintrunc } n \ (-1) = -1$
<proof>

lemma *bintrunc-Suc-numeral*:
 $\text{bintrunc } (\text{Suc } n) \ 1 = 1$
 $\text{bintrunc } (\text{Suc } n) \ (-1) = \text{bintrunc } n \ (-1) \text{ BIT } \text{True}$
 $\text{bintrunc } (\text{Suc } n) \ (\text{numeral } (\text{Num.Bit0 } w)) = \text{bintrunc } n \ (\text{numeral } w) \text{ BIT } \text{False}$
 $\text{bintrunc } (\text{Suc } n) \ (\text{numeral } (\text{Num.Bit1 } w)) = \text{bintrunc } n \ (\text{numeral } w) \text{ BIT } \text{True}$
 $\text{bintrunc } (\text{Suc } n) \ (- \text{numeral } (\text{Num.Bit0 } w)) =$
 $\quad \text{bintrunc } n \ (- \text{numeral } w) \text{ BIT } \text{False}$
 $\text{bintrunc } (\text{Suc } n) \ (- \text{numeral } (\text{Num.Bit1 } w)) =$
 $\quad \text{bintrunc } n \ (- \text{numeral } (w + \text{Num.One})) \text{ BIT } \text{True}$
<proof>

lemma *sbintrunc-0-numeral* [simp]:
 $\text{sbintrunc } 0 \ 1 = -1$
 $\text{sbintrunc } 0 \ (\text{numeral } (\text{Num.Bit0 } w)) = 0$
 $\text{sbintrunc } 0 \ (\text{numeral } (\text{Num.Bit1 } w)) = -1$

$sbintrunc\ 0\ (-\ numeral\ (Num.Bit0\ w)) = 0$
 $sbintrunc\ 0\ (-\ numeral\ (Num.Bit1\ w)) = -1$
 ⟨proof⟩

lemma *sbintrunc-Suc-numeral*:

$sbintrunc\ (Suc\ n)\ 1 = 1$
 $sbintrunc\ (Suc\ n)\ (numeral\ (Num.Bit0\ w)) =$
 $\quad sbintrunc\ n\ (numeral\ w)\ BIT\ False$
 $sbintrunc\ (Suc\ n)\ (numeral\ (Num.Bit1\ w)) =$
 $\quad sbintrunc\ n\ (numeral\ w)\ BIT\ True$
 $sbintrunc\ (Suc\ n)\ (-\ numeral\ (Num.Bit0\ w)) =$
 $\quad sbintrunc\ n\ (-\ numeral\ w)\ BIT\ False$
 $sbintrunc\ (Suc\ n)\ (-\ numeral\ (Num.Bit1\ w)) =$
 $\quad sbintrunc\ n\ (-\ numeral\ (w + Num.One))\ BIT\ True$
 ⟨proof⟩

lemma *bin-sign-lem*: $(bin-sign\ (sbintrunc\ n\ bin) = -1) = bin-nth\ bin\ n$
 ⟨proof⟩

lemma *nth-bintr*: $bin-nth\ (bintrunc\ m\ w)\ n = (n < m \ \& \ bin-nth\ w\ n)$
 ⟨proof⟩

lemma *nth-sbintr*:

$bin-nth\ (sbintrunc\ m\ w)\ n =$
 $\quad (if\ n < m\ then\ bin-nth\ w\ n\ else\ bin-nth\ w\ m)$
 ⟨proof⟩

lemma *bin-nth-Bit*:

$bin-nth\ (w\ BIT\ b)\ n = (n = 0 \ \& \ b \mid (EX\ m.\ n = Suc\ m \ \& \ bin-nth\ w\ m))$
 ⟨proof⟩

lemma *bin-nth-Bit0*:

$bin-nth\ (numeral\ (Num.Bit0\ w))\ n \longleftrightarrow$
 $\quad (\exists\ m.\ n = Suc\ m \ \wedge \ bin-nth\ (numeral\ w)\ m)$
 ⟨proof⟩

lemma *bin-nth-Bit1*:

$bin-nth\ (numeral\ (Num.Bit1\ w))\ n \longleftrightarrow$
 $\quad n = 0 \ \vee \ (\exists\ m.\ n = Suc\ m \ \wedge \ bin-nth\ (numeral\ w)\ m)$
 ⟨proof⟩

lemma *bintrunc-bintrunc-l*:

$n \leq m \implies (bintrunc\ m\ (bintrunc\ n\ w)) = bintrunc\ n\ w$
 ⟨proof⟩

lemma *sbintrunc-sbintrunc-l*:

$n \leq m \implies (sbintrunc\ m\ (sbintrunc\ n\ w)) = sbintrunc\ n\ w$
 ⟨proof⟩

lemma *bintrunc-bintrunc-ge*:

$$n \leq m \implies (\text{bintrunc } n (\text{bintrunc } m w) = \text{bintrunc } n w)$$

<proof>

lemma *bintrunc-bintrunc-min* [*simp*]:

$$\text{bintrunc } m (\text{bintrunc } n w) = \text{bintrunc } (\min m n) w$$

<proof>

lemma *sbintrunc-sbintrunc-min* [*simp*]:

$$\text{sbintrunc } m (\text{sbintrunc } n w) = \text{sbintrunc } (\min m n) w$$

<proof>

lemmas *bintrunc-Pls* =

$$\text{bintrunc.Suc} [\mathbf{where} \text{ bin}=0, \text{ simplified bin-last-numeral-simps bin-rest-numeral-simps}]$$

lemmas *bintrunc-Min* [*simp*] =

$$\text{bintrunc.Suc} [\mathbf{where} \text{ bin}=-1, \text{ simplified bin-last-numeral-simps bin-rest-numeral-simps}]$$

lemmas *bintrunc-BIT* [*simp*] =

$$\text{bintrunc.Suc} [\mathbf{where} \text{ bin}=w \text{ BIT } b, \text{ simplified bin-last-BIT bin-rest-BIT}] \mathbf{for } w b$$

lemmas *bintrunc-Sucs* = *bintrunc-Pls bintrunc-Min bintrunc-BIT*

bintrunc-Suc-numeral

lemmas *sbintrunc-Suc-Pls* =

$$\text{sbintrunc.Suc} [\mathbf{where} \text{ bin}=0, \text{ simplified bin-last-numeral-simps bin-rest-numeral-simps}]$$

lemmas *sbintrunc-Suc-Min* =

$$\text{sbintrunc.Suc} [\mathbf{where} \text{ bin}=-1, \text{ simplified bin-last-numeral-simps bin-rest-numeral-simps}]$$

lemmas *sbintrunc-Suc-BIT* [*simp*] =

$$\text{sbintrunc.Suc} [\mathbf{where} \text{ bin}=w \text{ BIT } b, \text{ simplified bin-last-BIT bin-rest-BIT}] \mathbf{for } w b$$

lemmas *sbintrunc-Sucs* = *sbintrunc-Suc-Pls sbintrunc-Suc-Min sbintrunc-Suc-BIT*

sbintrunc-Suc-numeral

lemmas *sbintrunc-Pls* =

$$\text{sbintrunc.Z} [\mathbf{where} \text{ bin}=0, \text{ simplified bin-last-numeral-simps bin-rest-numeral-simps}]$$

lemmas *sbintrunc-Min* =

$$\text{sbintrunc.Z} [\mathbf{where} \text{ bin}=-1, \text{ simplified bin-last-numeral-simps bin-rest-numeral-simps}]$$

lemmas *sbintrunc-0-BIT-B0* [*simp*] =

$$\text{sbintrunc.Z} [\mathbf{where} \text{ bin}=w \text{ BIT } \text{False}, \text{ simplified bin-last-numeral-simps bin-rest-numeral-simps}] \mathbf{for } w$$

lemmas *sbintrunc-0-BIT-B1* [*simp*] =
sbintrunc.Z [**where** *bin=w BIT True*,
simplified bin-last-BIT bin-rest-numeral-simps] **for** *w*

lemmas *sbintrunc-0-simps* =
sbintrunc-Pls sbintrunc-Min sbintrunc-0-BIT-B0 sbintrunc-0-BIT-B1

lemmas *bintrunc-simps* = *bintrunc.Z bintrunc-Sucs*
lemmas *sbintrunc-simps* = *sbintrunc-0-simps sbintrunc-Sucs*

lemma *bintrunc-minus*:
 $0 < n \implies \text{bintrunc } (\text{Suc } (n - 1)) w = \text{bintrunc } n w$
<proof>

lemma *sbintrunc-minus*:
 $0 < n \implies \text{sbintrunc } (\text{Suc } (n - 1)) w = \text{sbintrunc } n w$
<proof>

lemmas *bintrunc-minus-simps* =
bintrunc-Sucs [*THEN* [2] *bintrunc-minus* [*symmetric*, *THEN trans*]]
lemmas *sbintrunc-minus-simps* =
sbintrunc-Sucs [*THEN* [2] *sbintrunc-minus* [*symmetric*, *THEN trans*]]

lemmas *thobini1* = *arg-cong* [**where** *f = %w. w BIT b*] **for** *b*

lemmas *bintrunc-BIT-I* = *trans* [*OF bintrunc-BIT thobini1*]
lemmas *bintrunc-Min-I* = *trans* [*OF bintrunc-Min thobini1*]

lemmas *bmsts* = *bintrunc-minus-simps(1-3)* [*THEN thobini1* [*THEN* [2] *trans*]]
lemmas *bintrunc-Pls-minus-I* = *bmsts(1)*
lemmas *bintrunc-Min-minus-I* = *bmsts(2)*
lemmas *bintrunc-BIT-minus-I* = *bmsts(3)*

lemma *bintrunc-Suc-lem*:
 $\text{bintrunc } (\text{Suc } n) x = y \implies m = \text{Suc } n \implies \text{bintrunc } m x = y$
<proof>

lemmas *bintrunc-Suc-Ialts* =
bintrunc-Min-I [*THEN bintrunc-Suc-lem*]
bintrunc-BIT-I [*THEN bintrunc-Suc-lem*]

lemmas *sbintrunc-BIT-I* = *trans* [*OF sbintrunc-Suc-BIT thobini1*]

lemmas *sbintrunc-Suc-Is* =
sbintrunc-Sucs(1-3) [*THEN thobini1* [*THEN* [2] *trans*]]

lemmas *sbintrunc-Suc-minus-Is* =
sbintrunc-minus-simps(1-3) [*THEN thobini1* [*THEN* [2] *trans*]]

lemma *sbintrunc-Suc-lem*:

$$\text{sbintrunc } (\text{Suc } n) \ x = y \implies m = \text{Suc } n \implies \text{sbintrunc } m \ x = y$$

<proof>

lemmas *sbintrunc-Suc-Is* =

$$\text{sbintrunc-Suc-Is } [\text{THEN } \text{sbintrunc-Suc-lem}]$$

lemma *sbintrunc-bintrunc-lt*:

$$m > n \implies \text{sbintrunc } n \ (\text{bintrunc } m \ w) = \text{sbintrunc } n \ w$$

<proof>

lemma *bintrunc-sbintrunc-le*:

$$m \leq \text{Suc } n \implies \text{bintrunc } m \ (\text{sbintrunc } n \ w) = \text{bintrunc } m \ w$$

<proof>

lemmas *bintrunc-sbintrunc* [simp] = *order-refl* [THEN *bintrunc-sbintrunc-le*]

lemmas *sbintrunc-bintrunc* [simp] = *lessI* [THEN *sbintrunc-bintrunc-lt*]

lemmas *bintrunc-bintrunc* [simp] = *order-refl* [THEN *bintrunc-bintrunc-l*]

lemmas *sbintrunc-sbintrunc* [simp] = *order-refl* [THEN *sbintrunc-sbintrunc-l*]

lemma *bintrunc-sbintrunc'* [simp]:

$$0 < n \implies \text{bintrunc } n \ (\text{sbintrunc } (n - 1) \ w) = \text{bintrunc } n \ w$$

<proof>

lemma *sbintrunc-bintrunc'* [simp]:

$$0 < n \implies \text{sbintrunc } (n - 1) \ (\text{bintrunc } n \ w) = \text{sbintrunc } (n - 1) \ w$$

<proof>

lemma *bin-sbin-eq-iff*:

$$\text{bintrunc } (\text{Suc } n) \ x = \text{bintrunc } (\text{Suc } n) \ y \longleftrightarrow$$

$$\text{sbintrunc } n \ x = \text{sbintrunc } n \ y$$

<proof>

lemma *bin-sbin-eq-iff'*:

$$0 < n \implies \text{bintrunc } n \ x = \text{bintrunc } n \ y \longleftrightarrow$$

$$\text{sbintrunc } (n - 1) \ x = \text{sbintrunc } (n - 1) \ y$$

<proof>

lemmas *bintrunc-sbintruncS0* [simp] = *bintrunc-sbintrunc'* [unfolding *One-nat-def*]

lemmas *sbintrunc-bintruncS0* [simp] = *sbintrunc-bintrunc'* [unfolding *One-nat-def*]

lemmas *bintrunc-bintrunc-l'* = *le-add1* [THEN *bintrunc-bintrunc-l*]

lemmas *sbintrunc-sbintrunc-l'* = *le-add1* [THEN *sbintrunc-sbintrunc-l*]

lemmas *nat-non0-gr* =

$$\text{trans } [\text{OF } \text{iszero-def } [\text{THEN } \text{Not-eq-iff } [\text{THEN } \text{iffD2}]] \ \text{refl}]$$

lemma *bintrunc-numeral*:

bintrunc (numeral *k*) *x* =
bintrunc (pred-numeral *k*) (bin-rest *x*) BIT bin-last *x*
 ⟨proof⟩

lemma *sbintrunc-numeral*:

sbintrunc (numeral *k*) *x* =
sbintrunc (pred-numeral *k*) (bin-rest *x*) BIT bin-last *x*
 ⟨proof⟩

lemma *bintrunc-numeral-simps* [simp]:

bintrunc (numeral *k*) (numeral (Num.Bit0 *w*)) =
bintrunc (pred-numeral *k*) (numeral *w*) BIT False
bintrunc (numeral *k*) (numeral (Num.Bit1 *w*)) =
bintrunc (pred-numeral *k*) (numeral *w*) BIT True
bintrunc (numeral *k*) (– numeral (Num.Bit0 *w*)) =
bintrunc (pred-numeral *k*) (– numeral *w*) BIT False
bintrunc (numeral *k*) (– numeral (Num.Bit1 *w*)) =
bintrunc (pred-numeral *k*) (– numeral (*w* + Num.One)) BIT True
bintrunc (numeral *k*) 1 = 1
 ⟨proof⟩

lemma *sbintrunc-numeral-simps* [simp]:

sbintrunc (numeral *k*) (numeral (Num.Bit0 *w*)) =
sbintrunc (pred-numeral *k*) (numeral *w*) BIT False
sbintrunc (numeral *k*) (numeral (Num.Bit1 *w*)) =
sbintrunc (pred-numeral *k*) (numeral *w*) BIT True
sbintrunc (numeral *k*) (– numeral (Num.Bit0 *w*)) =
sbintrunc (pred-numeral *k*) (– numeral *w*) BIT False
sbintrunc (numeral *k*) (– numeral (Num.Bit1 *w*)) =
sbintrunc (pred-numeral *k*) (– numeral (*w* + Num.One)) BIT True
sbintrunc (numeral *k*) 1 = 1
 ⟨proof⟩

lemma *no-bintr-alt1*: *bintrunc* *n* = (λ*w*. *w* mod 2 ^ *n* :: int)

⟨proof⟩

lemma *range-bintrunc*: range (bintrunc *n*) = {*i*. 0 <= *i* & *i* < 2 ^ *n*}

⟨proof⟩

lemma *no-sbintr-alt2*:

sbintrunc *n* = (%*w*. (*w* + 2 ^ *n*) mod 2 ^ Suc *n* – 2 ^ *n* :: int)
 ⟨proof⟩

lemma *range-sbintrunc*:

range (sbintrunc *n*) = {*i*. – (2 ^ *n*) <= *i* & *i* < 2 ^ *n*}

⟨proof⟩

lemma *sb-inc-lem*:

$(a::int) + 2^k < 0 \implies a + 2^k + 2^{(Suc\ k)} \leq (a + 2^k) \bmod 2^{(Suc\ k)}$
 ⟨proof⟩

lemma *sb-inc-lem'*:

$(a::int) < -(2^k) \implies a + 2^k + 2^{(Suc\ k)} \leq (a + 2^k) \bmod 2^{(Suc\ k)}$
 ⟨proof⟩

lemma *sbintrunc-inc*:

$x < -(2^n) \implies x + 2^{(Suc\ n)} \leq \text{sbintrunc } n\ x$
 ⟨proof⟩

lemma *sb-dec-lem*:

$(0::int) \leq -(2^k) + a \implies (a + 2^k) \bmod (2 * 2^k) \leq -(2^k) + a$
 ⟨proof⟩

lemma *sb-dec-lem'*:

$(2::int)^k \leq a \implies (a + 2^k) \bmod (2 * 2^k) \leq -(2^k) + a$
 ⟨proof⟩

lemma *sbintrunc-dec*:

$x \geq (2^n) \implies x - 2^{(Suc\ n)} \geq \text{sbintrunc } n\ x$
 ⟨proof⟩

lemmas *zmod-uminus'* = *zminus-zmod* [where $m=c$] for c

lemmas *zpower-zmod'* = *power-mod* [where $b=c$ and $n=k$] for $c\ k$

lemmas *brdmod1s'* [symmetric] =

mod-add-left-eq mod-add-right-eq
mod-diff-left-eq mod-diff-right-eq
mod-mult-left-eq mod-mult-right-eq

lemmas *brdmods'* [symmetric] =

zpower-zmod' [symmetric]
trans [OF mod-add-left-eq mod-add-right-eq]
trans [OF mod-diff-left-eq mod-diff-right-eq]
trans [OF mod-mult-right-eq mod-mult-left-eq]
zmod-uminus' [symmetric]
mod-add-left-eq [where $b = 1::int$]
mod-diff-left-eq [where $b = 1::int$]

lemmas *bintr-arith1s* =

brdmod1s' [where $c=2^n::int$, folded *bintrunc-mod2p*] for n

lemmas *bintr-ariths* =

brdmods' [where $c=2^n::int$, folded *bintrunc-mod2p*] for n

lemmas *m2pths* = *pos-mod-sign pos-mod-bound* [OF *zless2p*]

lemma *bintr-ge0*: $0 \leq \text{bintrunc } n\ w$

⟨proof⟩

lemma *bintr-lt2p*: $\text{bintrunc } n \ w < 2 \wedge n$
 ⟨proof⟩

lemma *bintr-Min*: $\text{bintrunc } n \ (-1) = 2 \wedge n - 1$
 ⟨proof⟩

lemma *sbintr-ge*: $-(2 \wedge n) \leq \text{sbintrunc } n \ w$
 ⟨proof⟩

lemma *sbintr-lt*: $\text{sbintrunc } n \ w < 2 \wedge n$
 ⟨proof⟩

lemma *sign-Pls-ge-0*:
 $(\text{bin-sign } \text{bin} = 0) = (\text{bin} \geq (0 :: \text{int}))$
 ⟨proof⟩

lemma *sign-Min-lt-0*:
 $(\text{bin-sign } \text{bin} = -1) = (\text{bin} < (0 :: \text{int}))$
 ⟨proof⟩

lemma *bin-rest-trunc*:
 $(\text{bin-rest } (\text{bintrunc } n \ \text{bin})) = \text{bintrunc } (n - 1) (\text{bin-rest } \text{bin})$
 ⟨proof⟩

lemma *bin-rest-power-trunc*:
 $(\text{bin-rest } \wedge \wedge k) (\text{bintrunc } n \ \text{bin}) =$
 $\text{bintrunc } (n - k) ((\text{bin-rest } \wedge \wedge k) \ \text{bin})$
 ⟨proof⟩

lemma *bin-rest-trunc-i*:
 $\text{bintrunc } n \ (\text{bin-rest } \text{bin}) = \text{bin-rest } (\text{bintrunc } (\text{Suc } n) \ \text{bin})$
 ⟨proof⟩

lemma *bin-rest-strunc*:
 $\text{bin-rest } (\text{sbintrunc } (\text{Suc } n) \ \text{bin}) = \text{sbintrunc } n \ (\text{bin-rest } \text{bin})$
 ⟨proof⟩

lemma *bintrunc-rest [simp]*:
 $\text{bintrunc } n \ (\text{bin-rest } (\text{bintrunc } n \ \text{bin})) = \text{bin-rest } (\text{bintrunc } n \ \text{bin})$
 ⟨proof⟩

lemma *sbintrunc-rest [simp]*:
 $\text{sbintrunc } n \ (\text{bin-rest } (\text{sbintrunc } n \ \text{bin})) = \text{bin-rest } (\text{sbintrunc } n \ \text{bin})$
 ⟨proof⟩

lemma *bintrunc-rest'*:
 $\text{bintrunc } n \ o \ \text{bin-rest } o \ \text{bintrunc } n = \text{bin-rest } o \ \text{bintrunc } n$
 ⟨proof⟩

lemma *sbintrunc-rest'* :

sbintrunc n o bin-rest o sbintrunc n = bin-rest o sbintrunc n
 ⟨proof⟩

lemma *rco-lem*:

f o g o f = g o f ==> f o (g o f) ^^ n = g ^^ n o f
 ⟨proof⟩

lemmas *rco-bintr = bintrunc-rest'*

[*THEN rco-lem [THEN fun-cong], unfolded o-def*]

lemmas *rco-sbintr = sbintrunc-rest'*

[*THEN rco-lem [THEN fun-cong], unfolded o-def*]

10.4 Splitting and concatenation

primrec *bin-split* :: *nat* ⇒ *int* ⇒ *int* × *int* **where**

Z: *bin-split 0 w = (w, 0)*

| *Suc*: *bin-split (Suc n) w = (let (w1, w2) = bin-split n (bin-rest w)*
in (w1, w2 BIT bin-last w))

lemma [*code*]:

bin-split (Suc n) w = (let (w1, w2) = bin-split n (bin-rest w) in (w1, w2 BIT
bin-last w))

bin-split 0 w = (w, 0)

⟨proof⟩

primrec *bin-cat* :: *int* ⇒ *nat* ⇒ *int* ⇒ *int* **where**

Z: *bin-cat w 0 v = w*

| *Suc*: *bin-cat w (Suc n) v = bin-cat w n (bin-rest v) BIT bin-last v*

end

11 Bitwise Operations on Binary Integers

theory *Bits-Int*

imports *Bits Bit-Representation*

begin

11.1 Logical operations

bit-wise logical operations on the int type

instantiation *int* :: *bit*

begin

definition *int-not-def*:

bitNOT = (λx::int. - x - 1)

function *bitAND-int* **where**
bitAND-int x y =
 (if $x = 0$ then 0 else if $x = -1$ then y else
 (bin-rest x AND bin-rest y) BIT (bin-last x \wedge bin-last y))
 ⟨proof⟩

termination
 ⟨proof⟩

declare *bitAND-int.simps* [*simp del*]

definition *int-or-def*:
bitOR = ($\lambda x y::int. NOT (NOT x AND NOT y)$)

definition *int-xor-def*:
bitXOR = ($\lambda x y::int. (x AND NOT y) OR (NOT x AND y)$)

instance ⟨proof⟩

end

11.1.1 Basic simplification rules

lemma *int-not-BIT* [*simp*]:
 $NOT (w BIT b) = (NOT w) BIT (\neg b)$
 ⟨proof⟩

lemma *int-not-simps* [*simp*]:
 $NOT (0::int) = -1$
 $NOT (1::int) = -2$
 $NOT (-1::int) = 0$
 $NOT (numeral w::int) = - numeral (w + Num.One)$
 $NOT (- numeral (Num.Bit0 w)::int) = numeral (Num.BitM w)$
 $NOT (- numeral (Num.Bit1 w)::int) = numeral (Num.Bit0 w)$
 ⟨proof⟩

lemma *int-not-not* [*simp*]: $NOT (NOT (x::int)) = x$
 ⟨proof⟩

lemma *int-and-0* [*simp*]: $(0::int) AND x = 0$
 ⟨proof⟩

lemma *int-and-m1* [*simp*]: $(-1::int) AND x = x$
 ⟨proof⟩

lemma *int-and-Bits* [*simp*]:
 $(x BIT b) AND (y BIT c) = (x AND y) BIT (b \wedge c)$
 ⟨proof⟩

lemma *int-or-zero* [simp]: $(0::int) \text{ OR } x = x$
 ⟨proof⟩

lemma *int-or-minus1* [simp]: $(-1::int) \text{ OR } x = -1$
 ⟨proof⟩

lemma *int-or-Bits* [simp]:
 $(x \text{ BIT } b) \text{ OR } (y \text{ BIT } c) = (x \text{ OR } y) \text{ BIT } (b \vee c)$
 ⟨proof⟩

lemma *int-xor-zero* [simp]: $(0::int) \text{ XOR } x = x$
 ⟨proof⟩

lemma *int-xor-Bits* [simp]:
 $(x \text{ BIT } b) \text{ XOR } (y \text{ BIT } c) = (x \text{ XOR } y) \text{ BIT } ((b \vee c) \wedge \neg (b \wedge c))$
 ⟨proof⟩

11.1.2 Binary destructors

lemma *bin-rest-NOT* [simp]: $\text{bin-rest } (\text{NOT } x) = \text{NOT } (\text{bin-rest } x)$
 ⟨proof⟩

lemma *bin-last-NOT* [simp]: $\text{bin-last } (\text{NOT } x) \longleftrightarrow \neg \text{bin-last } x$
 ⟨proof⟩

lemma *bin-rest-AND* [simp]: $\text{bin-rest } (x \text{ AND } y) = \text{bin-rest } x \text{ AND } \text{bin-rest } y$
 ⟨proof⟩

lemma *bin-last-AND* [simp]: $\text{bin-last } (x \text{ AND } y) \longleftrightarrow \text{bin-last } x \wedge \text{bin-last } y$
 ⟨proof⟩

lemma *bin-rest-OR* [simp]: $\text{bin-rest } (x \text{ OR } y) = \text{bin-rest } x \text{ OR } \text{bin-rest } y$
 ⟨proof⟩

lemma *bin-last-OR* [simp]: $\text{bin-last } (x \text{ OR } y) \longleftrightarrow \text{bin-last } x \vee \text{bin-last } y$
 ⟨proof⟩

lemma *bin-rest-XOR* [simp]: $\text{bin-rest } (x \text{ XOR } y) = \text{bin-rest } x \text{ XOR } \text{bin-rest } y$
 ⟨proof⟩

lemma *bin-last-XOR* [simp]: $\text{bin-last } (x \text{ XOR } y) \longleftrightarrow (\text{bin-last } x \vee \text{bin-last } y) \wedge \neg (\text{bin-last } x \wedge \text{bin-last } y)$
 ⟨proof⟩

lemma *bin-nth-ops*:

!!x y. $\text{bin-nth } (x \text{ AND } y) \ n = (\text{bin-nth } x \ n \ \& \ \text{bin-nth } y \ n)$
 !!x y. $\text{bin-nth } (x \text{ OR } y) \ n = (\text{bin-nth } x \ n \ | \ \text{bin-nth } y \ n)$
 !!x y. $\text{bin-nth } (x \text{ XOR } y) \ n = (\text{bin-nth } x \ n \ \sim \ \text{bin-nth } y \ n)$
 !!x. $\text{bin-nth } (\text{NOT } x) \ n = (\sim \ \text{bin-nth } x \ n)$

<proof>

11.1.3 Derived properties

lemma *int-xor-minus1* [simp]: $(-1::int) \text{ XOR } x = \text{NOT } x$
<proof>

lemma *int-xor-extra-simps* [simp]:
 $w \text{ XOR } (0::int) = w$
 $w \text{ XOR } (-1::int) = \text{NOT } w$
<proof>

lemma *int-or-extra-simps* [simp]:
 $w \text{ OR } (0::int) = w$
 $w \text{ OR } (-1::int) = -1$
<proof>

lemma *int-and-extra-simps* [simp]:
 $w \text{ AND } (0::int) = 0$
 $w \text{ AND } (-1::int) = w$
<proof>

lemma *bin-ops-comm*:
shows
int-and-comm: $!!y::int. x \text{ AND } y = y \text{ AND } x$ **and**
int-or-comm: $!!y::int. x \text{ OR } y = y \text{ OR } x$ **and**
int-xor-comm: $!!y::int. x \text{ XOR } y = y \text{ XOR } x$
<proof>

lemma *bin-ops-same* [simp]:
 $(x::int) \text{ AND } x = x$
 $(x::int) \text{ OR } x = x$
 $(x::int) \text{ XOR } x = 0$
<proof>

lemmas *bin-log-esimps* =
int-and-extra-simps int-or-extra-simps int-xor-extra-simps
int-and-0 int-and-m1 int-or-zero int-or-minus1 int-xor-zero int-xor-minus1

lemma *bbw-ao-absorb*:
 $!!y::int. x \text{ AND } (y \text{ OR } x) = x \ \& \ x \text{ OR } (y \text{ AND } x) = x$
<proof>

lemma *bbw-ao-absorbs-other*:
 $x \text{ AND } (x \text{ OR } y) = x \ \wedge \ (y \text{ AND } x) \text{ OR } x = (x::int)$
 $(y \text{ OR } x) \text{ AND } x = x \ \wedge \ x \text{ OR } (x \text{ AND } y) = (x::int)$

$(x \text{ OR } y) \text{ AND } x = x \wedge (x \text{ AND } y) \text{ OR } x = (x::\text{int})$
 ⟨proof⟩

lemmas *bbw-ao-absorbs* [simp] = *bbw-ao-absorb bbw-ao-absorbs-other*

lemma *int-xor-not*:

!!y::int. (NOT x) XOR y = NOT (x XOR y) &
 x XOR (NOT y) = NOT (x XOR y)
 ⟨proof⟩

lemma *int-and-assoc*:

$(x \text{ AND } y) \text{ AND } (z::\text{int}) = x \text{ AND } (y \text{ AND } z)$
 ⟨proof⟩

lemma *int-or-assoc*:

$(x \text{ OR } y) \text{ OR } (z::\text{int}) = x \text{ OR } (y \text{ OR } z)$
 ⟨proof⟩

lemma *int-xor-assoc*:

$(x \text{ XOR } y) \text{ XOR } (z::\text{int}) = x \text{ XOR } (y \text{ XOR } z)$
 ⟨proof⟩

lemmas *bbw-assocs* = *int-and-assoc int-or-assoc int-xor-assoc*

lemma *bbw-lcs* [simp]:

$(y::\text{int}) \text{ AND } (x \text{ AND } z) = x \text{ AND } (y \text{ AND } z)$
 $(y::\text{int}) \text{ OR } (x \text{ OR } z) = x \text{ OR } (y \text{ OR } z)$
 $(y::\text{int}) \text{ XOR } (x \text{ XOR } z) = x \text{ XOR } (y \text{ XOR } z)$
 ⟨proof⟩

lemma *bbw-not-dist*:

!!y::int. NOT (x OR y) = (NOT x) AND (NOT y)
 !!y::int. NOT (x AND y) = (NOT x) OR (NOT y)
 ⟨proof⟩

lemma *bbw-oa-dist*:

!!y z::int. (x AND y) OR z =
 (x OR z) AND (y OR z)
 ⟨proof⟩

lemma *bbw-ao-dist*:

!!y z::int. (x OR y) AND z =
 (x AND z) OR (y AND z)
 ⟨proof⟩

11.1.4 Simplification with numerals

Cases for 0 and -1 are already covered by other simp rules.

lemma *bin-rl-eqI*: $\llbracket \text{bin-rest } x = \text{bin-rest } y; \text{bin-last } x = \text{bin-last } y \rrbracket \implies x = y$
 ⟨proof⟩

lemma *bin-rest-neg-numeral-BitM* [simp]:
 $\text{bin-rest } (- \text{numeral } (\text{Num.BitM } w)) = - \text{numeral } w$
 ⟨proof⟩

lemma *bin-last-neg-numeral-BitM* [simp]:
 $\text{bin-last } (- \text{numeral } (\text{Num.BitM } w))$
 ⟨proof⟩

FIXME: The rule sets below are very large (24 rules for each operator). Is there a simpler way to do this?

lemma *int-and-numerals* [simp]:
 $\text{numeral } (\text{Num.Bit0 } x) \text{ AND } \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ AND } \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ AND } \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ AND } \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ AND } \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ AND } \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ AND } \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ AND } \text{numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ AND } - \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ AND } - \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ AND } - \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ AND } - \text{numeral } (y + \text{Num.One})) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ AND } - \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ AND } - \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ AND } - \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ AND } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ AND } \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } x \text{ AND } \text{numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ AND } \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } x \text{ AND } \text{numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ AND } \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ AND } \text{numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ AND } \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ AND } \text{numeral } y) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ AND } - \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } x \text{ AND } - \text{numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ AND } - \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } x \text{ AND } - \text{numeral } (y + \text{Num.One})) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ AND } - \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ AND } - \text{numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ AND } - \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ AND } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $(1::\text{int}) \text{ AND } \text{numeral } (\text{Num.Bit0 } y) = 0$
 $(1::\text{int}) \text{ AND } \text{numeral } (\text{Num.Bit1 } y) = 1$
 $(1::\text{int}) \text{ AND } - \text{numeral } (\text{Num.Bit0 } y) = 0$

$(1::int) \text{ AND } - \text{ numeral } (\text{Num.Bit1 } y) = 1$
 $\text{ numeral } (\text{Num.Bit0 } x) \text{ AND } (1::int) = 0$
 $\text{ numeral } (\text{Num.Bit1 } x) \text{ AND } (1::int) = 1$
 $- \text{ numeral } (\text{Num.Bit0 } x) \text{ AND } (1::int) = 0$
 $- \text{ numeral } (\text{Num.Bit1 } x) \text{ AND } (1::int) = 1$
 ⟨proof⟩

lemma *int-or-numerals [simp]*:

$\text{ numeral } (\text{Num.Bit0 } x) \text{ OR numeral } (\text{Num.Bit0 } y) = (\text{ numeral } x \text{ OR numeral } y)$
BIT False
 $\text{ numeral } (\text{Num.Bit0 } x) \text{ OR numeral } (\text{Num.Bit1 } y) = (\text{ numeral } x \text{ OR numeral } y)$
BIT True
 $\text{ numeral } (\text{Num.Bit1 } x) \text{ OR numeral } (\text{Num.Bit0 } y) = (\text{ numeral } x \text{ OR numeral } y)$
BIT True
 $\text{ numeral } (\text{Num.Bit1 } x) \text{ OR numeral } (\text{Num.Bit1 } y) = (\text{ numeral } x \text{ OR numeral } y)$
BIT True
 $\text{ numeral } (\text{Num.Bit0 } x) \text{ OR } - \text{ numeral } (\text{Num.Bit0 } y) = (\text{ numeral } x \text{ OR } - \text{ numeral } y)$
BIT False
 $\text{ numeral } (\text{Num.Bit0 } x) \text{ OR } - \text{ numeral } (\text{Num.Bit1 } y) = (\text{ numeral } x \text{ OR } - \text{ numeral } (y + \text{Num.One}))$
BIT True
 $\text{ numeral } (\text{Num.Bit1 } x) \text{ OR } - \text{ numeral } (\text{Num.Bit0 } y) = (\text{ numeral } x \text{ OR } - \text{ numeral } y)$
BIT True
 $\text{ numeral } (\text{Num.Bit1 } x) \text{ OR } - \text{ numeral } (\text{Num.Bit1 } y) = (\text{ numeral } x \text{ OR } - \text{ numeral } (y + \text{Num.One}))$
BIT True
 $- \text{ numeral } (\text{Num.Bit0 } x) \text{ OR numeral } (\text{Num.Bit0 } y) = (- \text{ numeral } x \text{ OR numeral } y)$
BIT False
 $- \text{ numeral } (\text{Num.Bit0 } x) \text{ OR numeral } (\text{Num.Bit1 } y) = (- \text{ numeral } x \text{ OR numeral } y)$
BIT True
 $- \text{ numeral } (\text{Num.Bit1 } x) \text{ OR numeral } (\text{Num.Bit0 } y) = (- \text{ numeral } (x + \text{Num.One}) \text{ OR numeral } y)$
BIT True
 $- \text{ numeral } (\text{Num.Bit1 } x) \text{ OR numeral } (\text{Num.Bit1 } y) = (- \text{ numeral } (x + \text{Num.One}) \text{ OR numeral } y)$
BIT True
 $- \text{ numeral } (\text{Num.Bit0 } x) \text{ OR } - \text{ numeral } (\text{Num.Bit0 } y) = (- \text{ numeral } x \text{ OR } - \text{ numeral } y)$
BIT False
 $- \text{ numeral } (\text{Num.Bit0 } x) \text{ OR } - \text{ numeral } (\text{Num.Bit1 } y) = (- \text{ numeral } x \text{ OR } - \text{ numeral } (y + \text{Num.One}))$
BIT True
 $- \text{ numeral } (\text{Num.Bit1 } x) \text{ OR } - \text{ numeral } (\text{Num.Bit0 } y) = (- \text{ numeral } (x + \text{Num.One}) \text{ OR } - \text{ numeral } y)$
BIT True
 $- \text{ numeral } (\text{Num.Bit1 } x) \text{ OR } - \text{ numeral } (\text{Num.Bit1 } y) = (- \text{ numeral } (x + \text{Num.One}) \text{ OR } - \text{ numeral } (y + \text{Num.One}))$
BIT True
 $(1::int) \text{ OR numeral } (\text{Num.Bit0 } y) = \text{ numeral } (\text{Num.Bit1 } y)$
 $(1::int) \text{ OR numeral } (\text{Num.Bit1 } y) = \text{ numeral } (\text{Num.Bit1 } y)$
 $(1::int) \text{ OR } - \text{ numeral } (\text{Num.Bit0 } y) = - \text{ numeral } (\text{Num.BitM } y)$
 $(1::int) \text{ OR } - \text{ numeral } (\text{Num.Bit1 } y) = - \text{ numeral } (\text{Num.Bit1 } y)$
 $\text{ numeral } (\text{Num.Bit0 } x) \text{ OR } (1::int) = \text{ numeral } (\text{Num.Bit1 } x)$
 $\text{ numeral } (\text{Num.Bit1 } x) \text{ OR } (1::int) = \text{ numeral } (\text{Num.Bit1 } x)$
 $- \text{ numeral } (\text{Num.Bit0 } x) \text{ OR } (1::int) = - \text{ numeral } (\text{Num.BitM } x)$
 $- \text{ numeral } (\text{Num.Bit1 } x) \text{ OR } (1::int) = - \text{ numeral } (\text{Num.Bit1 } x)$
 ⟨proof⟩

lemma *int-xor-numerals* [simp]:

$\text{numeral } (\text{Num.Bit0 } x) \text{ XOR numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ XOR numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ XOR numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ XOR numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ XOR numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ XOR numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ XOR numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ XOR numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ XOR } - \text{numeral } y) \text{ BIT False}$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ XOR } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit0 } y) = (\text{numeral } x \text{ XOR } - \text{numeral } y) \text{ BIT True}$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit1 } y) = (\text{numeral } x \text{ XOR } - \text{numeral } (y + \text{Num.One})) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ XOR numeral } (\text{Num.Bit0 } y) = (- \text{numeral } x \text{ XOR numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ XOR numeral } (\text{Num.Bit1 } y) = (- \text{numeral } x \text{ XOR numeral } y) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ XOR numeral } (\text{Num.Bit0 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ XOR numeral } y) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ XOR numeral } (\text{Num.Bit1 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ XOR numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } x \text{ XOR } - \text{numeral } y) \text{ BIT False}$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } x \text{ XOR } - \text{numeral } (y + \text{Num.One})) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit0 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ XOR } - \text{numeral } y) \text{ BIT True}$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ XOR } - \text{numeral } (\text{Num.Bit1 } y) = (- \text{numeral } (x + \text{Num.One}) \text{ XOR } - \text{numeral } (y + \text{Num.One})) \text{ BIT False}$
 $(1::\text{int}) \text{ XOR numeral } (\text{Num.Bit0 } y) = \text{numeral } (\text{Num.Bit1 } y)$
 $(1::\text{int}) \text{ XOR numeral } (\text{Num.Bit1 } y) = \text{numeral } (\text{Num.Bit0 } y)$
 $(1::\text{int}) \text{ XOR } - \text{numeral } (\text{Num.Bit0 } y) = - \text{numeral } (\text{Num.BitM } y)$
 $(1::\text{int}) \text{ XOR } - \text{numeral } (\text{Num.Bit1 } y) = - \text{numeral } (\text{Num.Bit0 } (y + \text{Num.One}))$
 $\text{numeral } (\text{Num.Bit0 } x) \text{ XOR } (1::\text{int}) = \text{numeral } (\text{Num.Bit1 } x)$
 $\text{numeral } (\text{Num.Bit1 } x) \text{ XOR } (1::\text{int}) = \text{numeral } (\text{Num.Bit0 } x)$
 $- \text{numeral } (\text{Num.Bit0 } x) \text{ XOR } (1::\text{int}) = - \text{numeral } (\text{Num.BitM } x)$
 $- \text{numeral } (\text{Num.Bit1 } x) \text{ XOR } (1::\text{int}) = - \text{numeral } (\text{Num.Bit0 } (x + \text{Num.One}))$
 ⟨proof⟩

11.1.5 Interactions with arithmetic

lemma *plus-and-or* [rule-format]:

$\text{ALL } y::\text{int. } (x \text{ AND } y) + (x \text{ OR } y) = x + y$
 ⟨proof⟩

lemma *le-int-or*:

$bin\text{-}sign\ (y::int) = 0 \implies x \leq x\ OR\ y$
 $\langle proof \rangle$

lemmas *int-and-le* =

$xtrans(3)\ [OF\ bbw\text{-}ao\text{-}absorbs\ (2)\ [THEN\ conjunct2,\ symmetric]\ le\text{-}int\text{-}or]$

lemma *bin-add-not*: $x + NOT\ x = (-1::int)$

$\langle proof \rangle$

11.1.6 Truncating results of bit-wise operations

lemma *bin-trunc-ao*:

$!!x\ y.\ (bintrunc\ n\ x)\ AND\ (bintrunc\ n\ y) = bintrunc\ n\ (x\ AND\ y)$
 $!!x\ y.\ (bintrunc\ n\ x)\ OR\ (bintrunc\ n\ y) = bintrunc\ n\ (x\ OR\ y)$
 $\langle proof \rangle$

lemma *bin-trunc-xor*:

$!!x\ y.\ bintrunc\ n\ (bintrunc\ n\ x\ XOR\ bintrunc\ n\ y) =$
 $bintrunc\ n\ (x\ XOR\ y)$
 $\langle proof \rangle$

lemma *bin-trunc-not*:

$!!x.\ bintrunc\ n\ (NOT\ (bintrunc\ n\ x)) = bintrunc\ n\ (NOT\ x)$
 $\langle proof \rangle$

lemma *bintr-bintr-i*:

$x = bintrunc\ n\ y \implies bintrunc\ n\ x = bintrunc\ n\ y$
 $\langle proof \rangle$

lemmas *bin-trunc-and* = *bin-trunc-ao*(1) [THEN *bintr-bintr-i*]

lemmas *bin-trunc-or* = *bin-trunc-ao*(2) [THEN *bintr-bintr-i*]

11.2 Setting and clearing bits

primrec

$bin\text{-}sc :: nat \Rightarrow bool \Rightarrow int \Rightarrow int$

where

$Z:$ $bin\text{-}sc\ 0\ b\ w = bin\text{-}rest\ w\ BIT\ b$

$|$ $Suc:$ $bin\text{-}sc\ (Suc\ n)\ b\ w = bin\text{-}sc\ n\ b\ (bin\text{-}rest\ w)\ BIT\ bin\text{-}last\ w$

lemma *bin-nth-sc* [simp]:

$bin\text{-}nth\ (bin\text{-}sc\ n\ b\ w)\ n \longleftrightarrow b$
 $\langle proof \rangle$

lemma *bin-sc-sc-same* [simp]:

$bin-sc\ n\ c\ (bin-sc\ n\ b\ w) = bin-sc\ n\ c\ w$
 ⟨proof⟩

lemma *bin-sc-sc-diff*:

$m \sim = n \implies$
 $bin-sc\ m\ c\ (bin-sc\ n\ b\ w) = bin-sc\ n\ b\ (bin-sc\ m\ c\ w)$
 ⟨proof⟩

lemma *bin-nth-sc-gen*:

$bin-nth\ (bin-sc\ n\ b\ w)\ m = (if\ m = n\ then\ b\ else\ bin-nth\ w\ m)$
 ⟨proof⟩

lemma *bin-sc-nth [simp]*:

$(bin-sc\ n\ (bin-nth\ w\ n)\ w) = w$
 ⟨proof⟩

lemma *bin-sign-sc [simp]*:

$bin-sign\ (bin-sc\ n\ b\ w) = bin-sign\ w$
 ⟨proof⟩

lemma *bin-sc-bintr [simp]*:

$bintrunc\ m\ (bin-sc\ n\ x\ (bintrunc\ m\ (w))) = bintrunc\ m\ (bin-sc\ n\ x\ w)$
 ⟨proof⟩

lemma *bin-clr-le*:

$bin-sc\ n\ False\ w \leq w$
 ⟨proof⟩

lemma *bin-set-ge*:

$bin-sc\ n\ True\ w \geq w$
 ⟨proof⟩

lemma *bintr-bin-clr-le*:

$bintrunc\ n\ (bin-sc\ m\ False\ w) \leq bintrunc\ n\ w$
 ⟨proof⟩

lemma *bintr-bin-set-ge*:

$bintrunc\ n\ (bin-sc\ m\ True\ w) \geq bintrunc\ n\ w$
 ⟨proof⟩

lemma *bin-sc-FP [simp]*: $bin-sc\ n\ False\ 0 = 0$

⟨proof⟩

lemma *bin-sc-TM [simp]*: $bin-sc\ n\ True\ (-\ 1) = -\ 1$

⟨proof⟩

lemmas *bin-sc-simps = bin-sc.Z bin-sc.Suc bin-sc-TM bin-sc-FP*

lemma *bin-sc-minus*:

$0 < n \implies \text{bin-sc } (\text{Suc } (n - 1)) \text{ } b \text{ } w = \text{bin-sc } n \text{ } b \text{ } w$
 ⟨proof⟩

lemmas *bin-sc-Suc-minus* =
trans [OF bin-sc-minus [symmetric] bin-sc.Suc]

lemma *bin-sc-numeral* [simp]:
 $\text{bin-sc } (\text{numeral } k) \text{ } b \text{ } w =$
 $\text{bin-sc } (\text{pred-numeral } k) \text{ } b \text{ } (\text{bin-rest } w) \text{ } \text{BIT } \text{bin-last } w$
 ⟨proof⟩

11.3 Splitting and concatenation

definition *bin-rcat* :: $\text{nat} \Rightarrow \text{int list} \Rightarrow \text{int}$
where

$\text{bin-rcat } n = \text{foldl } (\lambda u \ v. \text{bin-cat } u \ n \ v) \ 0$

fun *bin-rsplit-aux* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{int} \Rightarrow \text{int list} \Rightarrow \text{int list}$
where

$\text{bin-rsplit-aux } n \ m \ c \ bs =$
 (if $m = 0 \mid n = 0$ then bs else
 let $(a, b) = \text{bin-split } n \ c$
 in $\text{bin-rsplit-aux } n \ (m - n) \ a \ (b \# \ bs)$)

definition *bin-rsplit* :: $\text{nat} \Rightarrow \text{nat} \times \text{int} \Rightarrow \text{int list}$
where

$\text{bin-rsplit } n \ w = \text{bin-rsplit-aux } n \ (\text{fst } w) \ (\text{snd } w) \ []$

fun *bin-rsplittl-aux* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{int} \Rightarrow \text{int list} \Rightarrow \text{int list}$
where

$\text{bin-rsplittl-aux } n \ m \ c \ bs =$
 (if $m = 0 \mid n = 0$ then bs else
 let $(a, b) = \text{bin-split } (\text{min } m \ n) \ c$
 in $\text{bin-rsplittl-aux } n \ (m - n) \ a \ (b \# \ bs)$)

definition *bin-rsplittl* :: $\text{nat} \Rightarrow \text{nat} \times \text{int} \Rightarrow \text{int list}$
where

$\text{bin-rsplittl } n \ w = \text{bin-rsplittl-aux } n \ (\text{fst } w) \ (\text{snd } w) \ []$

declare *bin-rsplit-aux.simps* [simp del]

declare *bin-rsplittl-aux.simps* [simp del]

lemma *bin-sign-cat*:

$\text{bin-sign } (\text{bin-cat } x \ n \ y) = \text{bin-sign } x$
 ⟨proof⟩

lemma *bin-cat-Suc-Bit*:

$\text{bin-cat } w \ (\text{Suc } n) \ (v \ \text{BIT } b) = \text{bin-cat } w \ n \ v \ \text{BIT } b$
 ⟨proof⟩

lemma *bin-nth-cat*:

$bin_nth (bin_cat\ x\ k\ y)\ n =$
 (if $n < k$ then $bin_nth\ y\ n$ else $bin_nth\ x\ (n - k)$)
 ⟨proof⟩

lemma *bin-nth-split*:

$bin_split\ n\ c = (a, b) ==>$
 (ALL k . $bin_nth\ a\ k = bin_nth\ c\ (n + k)$) &
 (ALL k . $bin_nth\ b\ k = (k < n \ \&\ bin_nth\ c\ k)$)
 ⟨proof⟩

lemma *bin-cat-assoc*:

$bin_cat (bin_cat\ x\ m\ y)\ n\ z = bin_cat\ x\ (m + n)\ (bin_cat\ y\ n\ z)$
 ⟨proof⟩

lemma *bin-cat-assoc-sym*:

$bin_cat\ x\ m\ (bin_cat\ y\ n\ z) = bin_cat (bin_cat\ x\ (m - n)\ y)\ (min\ m\ n)\ z$
 ⟨proof⟩

lemma *bin-cat-zero* [simp]: $bin_cat\ 0\ n\ w = bintrunc\ n\ w$

⟨proof⟩

lemma *bintr-cat1*:

$bintrunc\ (k + n)\ (bin_cat\ a\ n\ b) = bin_cat (bintrunc\ k\ a)\ n\ b$
 ⟨proof⟩

lemma *bintr-cat*: $bintrunc\ m\ (bin_cat\ a\ n\ b) =$

$bin_cat (bintrunc\ (m - n)\ a)\ n\ (bintrunc\ (min\ m\ n)\ b)$
 ⟨proof⟩

lemma *bintr-cat-same* [simp]:

$bintrunc\ n\ (bin_cat\ a\ n\ b) = bintrunc\ n\ b$
 ⟨proof⟩

lemma *cat-bintr* [simp]:

$bin_cat\ a\ n\ (bintrunc\ n\ b) = bin_cat\ a\ n\ b$
 ⟨proof⟩

lemma *split-bintrunc*:

$bin_split\ n\ c = (a, b) ==> b = bintrunc\ n\ c$
 ⟨proof⟩

lemma *bin-cat-split*:

$bin_split\ n\ w = (u, v) ==> w = bin_cat\ u\ n\ v$
 ⟨proof⟩

lemma *bin-split-cat*:

$bin_split\ n\ (bin_cat\ v\ n\ w) = (v, bintrunc\ n\ w)$

<proof>

lemma *bin-split-zero* [*simp*]: $\text{bin-split } n \ 0 = (0, 0)$
<proof>

lemma *bin-split-minus1* [*simp*]:
 $\text{bin-split } n \ (-1) = (-1, \text{bintrunc } n \ (-1))$
<proof>

lemma *bin-split-trunc*:
 $\text{bin-split } (\text{min } m \ n) \ c = (a, b) \implies$
 $\text{bin-split } n \ (\text{bintrunc } m \ c) = (\text{bintrunc } (m - n) \ a, b)$
<proof>

lemma *bin-split-trunc1*:
 $\text{bin-split } n \ c = (a, b) \implies$
 $\text{bin-split } n \ (\text{bintrunc } m \ c) = (\text{bintrunc } (m - n) \ a, \text{bintrunc } m \ b)$
<proof>

lemma *bin-cat-num*:
 $\text{bin-cat } a \ n \ b = a * 2 ^ n + \text{bintrunc } n \ b$
<proof>

lemma *bin-split-num*:
 $\text{bin-split } n \ b = (b \ \text{div } 2 ^ n, b \ \text{mod } 2 ^ n)$
<proof>

11.4 Miscellaneous lemmas

lemma *nth-2p-bin*:
 $\text{bin-nth } (2 ^ n) \ m = (m = n)$
<proof>

lemma *ex-eq-or*:
 $(\text{EX } m. \ n = \text{Suc } m \ \& \ (m = k \ | \ P \ m)) = (n = \text{Suc } k \ | \ (\text{EX } m. \ n = \text{Suc } m \ \& \ P \ m))$
<proof>

lemma *power-BIT*: $2 ^ (\text{Suc } n) - 1 = (2 ^ n - 1) \ \text{BIT } \text{True}$
<proof>

lemma *mod-BIT*:
 $\text{bin } \text{BIT } \text{bit } \ \text{mod } 2 ^ \text{Suc } n = (\text{bin } \ \text{mod } 2 ^ n) \ \text{BIT } \text{bit}$
<proof>

lemma *AND-mod*:
fixes $x :: \text{int}$

shows $x \text{ AND } 2^n - 1 = x \text{ mod } 2^n$
 ⟨proof⟩

end

12 Bool lists and integers

theory *Bool-List-Representation*

imports *Main Bits-Int*

begin

definition $\text{map2} :: ('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a \text{ list} \Rightarrow 'b \text{ list} \Rightarrow 'c \text{ list}$

where

$\text{map2 } f \text{ as } bs = \text{map } (\text{case-prod } f) (\text{zip as } bs)$

lemma map2-Nil [*simp, code*]:

$\text{map2 } f [] \text{ ys} = []$

⟨proof⟩

lemma map2-Nil2 [*simp, code*]:

$\text{map2 } f \text{ xs } [] = []$

⟨proof⟩

lemma map2-Cons [*simp, code*]:

$\text{map2 } f (x \# \text{xs}) (y \# \text{ys}) = f \ x \ y \# \text{map2 } f \ \text{xs} \ \text{ys}$

⟨proof⟩

12.1 Operations on lists of booleans

primrec $\text{bl-to-bin-aux} :: \text{bool list} \Rightarrow \text{int} \Rightarrow \text{int}$

where

Nil: $\text{bl-to-bin-aux } [] \ w = w$

| *Cons*: $\text{bl-to-bin-aux } (b \# \text{bs}) \ w =$
 $\text{bl-to-bin-aux } \text{bs} \ (w \text{ BIT } b)$

definition $\text{bl-to-bin} :: \text{bool list} \Rightarrow \text{int}$

where

bl-to-bin-def : $\text{bl-to-bin } \text{bs} = \text{bl-to-bin-aux } \text{bs} \ 0$

primrec $\text{bin-to-bl-aux} :: \text{nat} \Rightarrow \text{int} \Rightarrow \text{bool list} \Rightarrow \text{bool list}$

where

Z: $\text{bin-to-bl-aux } 0 \ w \ \text{bl} = \text{bl}$

| *Suc*: $\text{bin-to-bl-aux } (\text{Suc } n) \ w \ \text{bl} =$
 $\text{bin-to-bl-aux } n \ (\text{bin-rest } w) \ ((\text{bin-last } w) \# \text{bl})$

definition $\text{bin-to-bl} :: \text{nat} \Rightarrow \text{int} \Rightarrow \text{bool list}$

where

bin-to-bl-def : $\text{bin-to-bl } n \ w = \text{bin-to-bl-aux } n \ w \ []$

primrec *bl-of-nth* :: *nat* \Rightarrow (*nat* \Rightarrow *bool*) \Rightarrow *bool list*
where
Suc: *bl-of-nth* (*Suc n*) *f* = *f n* # *bl-of-nth n f*
| *Z*: *bl-of-nth 0 f* = []

primrec *takefill* :: '*a* \Rightarrow *nat* \Rightarrow '*a list* \Rightarrow '*a list*
where
Z: *takefill fill 0 xs* = []
| *Suc*: *takefill fill (Suc n) xs* = (
 case xs of [] \Rightarrow *fill* # *takefill fill n xs*
 | *y* # *ys* \Rightarrow *y* # *takefill fill n ys*)

12.2 Arithmetic in terms of bool lists

Arithmetic operations in terms of the reversed bool list, assuming input list(s) the same length, and don't extend them.

primrec *rbl-succ* :: *bool list* \Rightarrow *bool list*
where
Nil: *rbl-succ Nil* = *Nil*
| *Cons*: *rbl-succ (x # xs)* = (*if x then False* # *rbl-succ xs else True* # *xs*)

primrec *rbl-pred* :: *bool list* \Rightarrow *bool list*
where
Nil: *rbl-pred Nil* = *Nil*
| *Cons*: *rbl-pred (x # xs)* = (*if x then False* # *xs else True* # *rbl-pred xs*)

primrec *rbl-add* :: *bool list* \Rightarrow *bool list* \Rightarrow *bool list*
where
— result is length of first arg, second arg may be longer
Nil: *rbl-add Nil x* = *Nil*
| *Cons*: *rbl-add (y # ys) x* = (*let ws = rbl-add ys (tl x) in*
 (*y* \sim *hd x*) # (*if hd x & y then rbl-succ ws else ws*)

primrec *rbl-mult* :: *bool list* \Rightarrow *bool list* \Rightarrow *bool list*
where
— result is length of first arg, second arg may be longer
Nil: *rbl-mult Nil x* = *Nil*
| *Cons*: *rbl-mult (y # ys) x* = (*let ws = False* # *rbl-mult ys x in*
 if y then rbl-add ws x else ws)

lemma *butlast-power*:
(*butlast* $\hat{\hat{}}$ *n*) *bl* = *take (length bl - n) bl*
<proof>

lemma *bin-to-bl-aux-zero-minus-simp* [*simp*]:
 $0 < n \implies \text{bin-to-bl-aux } n \ 0 \ bl =$
 $\text{bin-to-bl-aux } (n - 1) \ 0 \ (\text{False} \ \# \ bl)$
<proof>

lemma *bin-to-bl-aux-minus1-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ (-1) \ \text{bl} = \\ \text{bin-to-bl-aux } (n - 1) \ (-1) \ (\text{True} \ \# \ \text{bl}) \\ \langle \text{proof} \rangle$$

lemma *bin-to-bl-aux-one-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ 1 \ \text{bl} = \\ \text{bin-to-bl-aux } (n - 1) \ 0 \ (\text{True} \ \# \ \text{bl}) \\ \langle \text{proof} \rangle$$

lemma *bin-to-bl-aux-Bit-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ (w \ \text{BIT} \ b) \ \text{bl} = \\ \text{bin-to-bl-aux } (n - 1) \ w \ (b \ \# \ \text{bl}) \\ \langle \text{proof} \rangle$$

lemma *bin-to-bl-aux-Bit0-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ (\text{numeral } (\text{Num.Bit0 } w)) \ \text{bl} = \\ \text{bin-to-bl-aux } (n - 1) \ (\text{numeral } w) \ (\text{False} \ \# \ \text{bl}) \\ \langle \text{proof} \rangle$$

lemma *bin-to-bl-aux-Bit1-minus-simp* [simp]:

$$0 < n \implies \text{bin-to-bl-aux } n \ (\text{numeral } (\text{Num.Bit1 } w)) \ \text{bl} = \\ \text{bin-to-bl-aux } (n - 1) \ (\text{numeral } w) \ (\text{True} \ \# \ \text{bl}) \\ \langle \text{proof} \rangle$$

Link between bin and bool list.

lemma *bl-to-bin-aux-append*:

$$\text{bl-to-bin-aux } (bs \ @ \ cs) \ w = \text{bl-to-bin-aux } cs \ (\text{bl-to-bin-aux } bs \ w) \\ \langle \text{proof} \rangle$$

lemma *bin-to-bl-aux-append*:

$$\text{bin-to-bl-aux } n \ w \ bs \ @ \ cs = \text{bin-to-bl-aux } n \ w \ (bs \ @ \ cs) \\ \langle \text{proof} \rangle$$

lemma *bl-to-bin-append*:

$$\text{bl-to-bin } (bs \ @ \ cs) = \text{bl-to-bin-aux } cs \ (\text{bl-to-bin } bs) \\ \langle \text{proof} \rangle$$

lemma *bin-to-bl-aux-alt*:

$$\text{bin-to-bl-aux } n \ w \ bs = \text{bin-to-bl } n \ w \ @ \ bs \\ \langle \text{proof} \rangle$$

lemma *bin-to-bl-0* [simp]: *bin-to-bl* 0 *bs* = []

$\langle \text{proof} \rangle$

lemma *size-bin-to-bl-aux*:

$$\text{size } (\text{bin-to-bl-aux } n \ w \ bs) = n + \text{length } bs \\ \langle \text{proof} \rangle$$

lemma *size-bin-to-bl* [simp]: $\text{size } (\text{bin-to-bl } n \ w) = n$
 ⟨proof⟩

lemma *bin-bl-bin'*:
 $\text{bl-to-bin } (\text{bin-to-bl-aux } n \ w \ bs) =$
 $\text{bl-to-bin-aux } bs \ (\text{bintrunc } n \ w)$
 ⟨proof⟩

lemma *bin-bl-bin* [simp]: $\text{bl-to-bin } (\text{bin-to-bl } n \ w) = \text{bintrunc } n \ w$
 ⟨proof⟩

lemma *bl-bin-bl'*:
 $\text{bin-to-bl } (n + \text{length } bs) \ (\text{bl-to-bin-aux } bs \ w) =$
 $\text{bin-to-bl-aux } n \ w \ bs$
 ⟨proof⟩

lemma *bl-bin-bl* [simp]: $\text{bin-to-bl } (\text{length } bs) \ (\text{bl-to-bin } bs) = bs$
 ⟨proof⟩

lemma *bl-to-bin-inj*:
 $\text{bl-to-bin } bs = \text{bl-to-bin } cs \implies \text{length } bs = \text{length } cs \implies bs = cs$
 ⟨proof⟩

lemma *bl-to-bin-False* [simp]: $\text{bl-to-bin } (\text{False} \ \# \ bl) = \text{bl-to-bin } bl$
 ⟨proof⟩

lemma *bl-to-bin-Nil* [simp]: $\text{bl-to-bin } [] = 0$
 ⟨proof⟩

lemma *bin-to-bl-zero-aux*:
 $\text{bin-to-bl-aux } n \ 0 \ bl = \text{replicate } n \ \text{False} \ @ \ bl$
 ⟨proof⟩

lemma *bin-to-bl-zero*: $\text{bin-to-bl } n \ 0 = \text{replicate } n \ \text{False}$
 ⟨proof⟩

lemma *bin-to-bl-minus1-aux*:
 $\text{bin-to-bl-aux } n \ (-1) \ bl = \text{replicate } n \ \text{True} \ @ \ bl$
 ⟨proof⟩

lemma *bin-to-bl-minus1*: $\text{bin-to-bl } n \ (-1) = \text{replicate } n \ \text{True}$
 ⟨proof⟩

lemma *bl-to-bin-rep-F*:
 $\text{bl-to-bin } (\text{replicate } n \ \text{False} \ @ \ bl) = \text{bl-to-bin } bl$
 ⟨proof⟩

lemma *bin-to-bl-trunc* [simp]:
 $n \leq m \implies \text{bin-to-bl } n \ (\text{bintrunc } m \ w) = \text{bin-to-bl } n \ w$

<proof>

lemma *bin-to-bl-aux-bintr*:

$bin-to-bl-aux\ n\ (bintrunc\ m\ bin)\ bl =$
 $replicate\ (n - m)\ False\ @\ bin-to-bl-aux\ (min\ n\ m)\ bin\ bl$
<proof>

lemma *bin-to-bl-bintr*:

$bin-to-bl\ n\ (bintrunc\ m\ bin) =$
 $replicate\ (n - m)\ False\ @\ bin-to-bl\ (min\ n\ m)\ bin$
<proof>

lemma *bl-to-bin-rep-False*: $bl-to-bin\ (replicate\ n\ False) = 0$

<proof>

lemma *len-bin-to-bl-aux*:

$length\ (bin-to-bl-aux\ n\ w\ bs) = n + length\ bs$
<proof>

lemma *len-bin-to-bl [simp]*: $length\ (bin-to-bl\ n\ w) = n$

<proof>

lemma *sign-bl-bin'*:

$bin-sign\ (bl-to-bin-aux\ bs\ w) = bin-sign\ w$
<proof>

lemma *sign-bl-bin*: $bin-sign\ (bl-to-bin\ bs) = 0$

<proof>

lemma *bl-sbin-sign-aux*:

$hd\ (bin-to-bl-aux\ (Suc\ n)\ w\ bs) =$
 $(bin-sign\ (sbintrunc\ n\ w) = -1)$
<proof>

lemma *bl-sbin-sign*:

$hd\ (bin-to-bl\ (Suc\ n)\ w) = (bin-sign\ (sbintrunc\ n\ w) = -1)$
<proof>

lemma *bin-nth-of-bl-aux*:

$bin-nth\ (bl-to-bin-aux\ bl\ w)\ n =$
 $(n < size\ bl \ \&\ rev\ bl\ !\ n \ | \ n \geq length\ bl \ \&\ bin-nth\ w\ (n - size\ bl))$
<proof>

lemma *bin-nth-of-bl*: $bin-nth\ (bl-to-bin\ bl)\ n = (n < length\ bl \ \&\ rev\ bl\ !\ n)$

<proof>

lemma *bin-nth-bl*: $n < m \implies bin-nth\ w\ n = nth\ (rev\ (bin-to-bl\ m\ w))\ n$

<proof>

lemma *nth-rev*:

$$n < \text{length } xs \implies \text{rev } xs ! n = xs ! (\text{length } xs - 1 - n)$$

<proof>

lemma *nth-rev-alt*: $n < \text{length } ys \implies ys ! n = \text{rev } ys ! (\text{length } ys - \text{Suc } n)$

<proof>

lemma *nth-bin-to-bl-aux*:

$$n < m + \text{length } bl \implies (\text{bin-to-bl-aux } m \ w \ bl) ! n = \\ (\text{if } n < m \text{ then } \text{bin-nth } w \ (m - 1 - n) \text{ else } bl ! (n - m))$$

<proof>

lemma *nth-bin-to-bl*: $n < m \implies (\text{bin-to-bl } m \ w) ! n = \text{bin-nth } w \ (m - \text{Suc } n)$

<proof>

lemma *bl-to-bin-lt2p-aux*:

$$\text{bl-to-bin-aux } bs \ w < (w + 1) * (2 \wedge \text{length } bs)$$

<proof>

lemma *bl-to-bin-lt2p-drop*:

$$\text{bl-to-bin } bs < 2 \wedge \text{length } (\text{dropWhile Not } bs)$$

<proof>

lemma *bl-to-bin-lt2p*: $\text{bl-to-bin } bs < 2 \wedge \text{length } bs$

<proof>

lemma *bl-to-bin-ge2p-aux*:

$$\text{bl-to-bin-aux } bs \ w \geq w * (2 \wedge \text{length } bs)$$

<proof>

lemma *bl-to-bin-ge0*: $\text{bl-to-bin } bs \geq 0$

<proof>

lemma *butlast-rest-bin*:

$$\text{butlast } (\text{bin-to-bl } n \ w) = \text{bin-to-bl } (n - 1) \ (\text{bin-rest } w)$$

<proof>

lemma *butlast-bin-rest*:

$$\text{butlast } bl = \text{bin-to-bl } (\text{length } bl - \text{Suc } 0) \ (\text{bin-rest } (\text{bl-to-bin } bl))$$

<proof>

lemma *butlast-rest-bl2bin-aux*:

$$bl \sim [] \implies$$

$$\text{bl-to-bin-aux } (\text{butlast } bl) \ w = \text{bin-rest } (\text{bl-to-bin-aux } bl \ w)$$

<proof>

lemma *butlast-rest-bl2bin*:

$$\text{bl-to-bin } (\text{butlast } bl) = \text{bin-rest } (\text{bl-to-bin } bl)$$

<proof>

lemma *trunc-bl2bin-aux*:

$$\text{bintrunc } m \text{ (bl-to-bin-aux bl w) =}$$

$$\text{bl-to-bin-aux (drop (length bl - m) bl) (bintrunc (m - length bl) w)}$$

⟨proof⟩

lemma *trunc-bl2bin*:

$$\text{bintrunc } m \text{ (bl-to-bin bl) = bl-to-bin (drop (length bl - m) bl)}$$

⟨proof⟩

lemma *trunc-bl2bin-len [simp]*:

$$\text{bintrunc (length bl) (bl-to-bin bl) = bl-to-bin bl}$$

⟨proof⟩

lemma *bl2bin-drop*:

$$\text{bl-to-bin (drop k bl) = bintrunc (length bl - k) (bl-to-bin bl)}$$

⟨proof⟩

lemma *nth-rest-power-bin*:

$$\text{bin-nth ((bin-rest ^ k) w) n = bin-nth w (n + k)}$$

⟨proof⟩

lemma *take-rest-power-bin*:

$$m \leq n \implies \text{take } m \text{ (bin-to-bl } n \text{ w) = bin-to-bl } m \text{ ((bin-rest ^ (n - m)) w)}$$

⟨proof⟩

lemma *hd-butlast*: $\text{size } xs > 1 \implies \text{hd (butlast xs) = hd xs}$

⟨proof⟩

lemma *last-bin-last'*:

$$\text{size } xs > 0 \implies \text{last } xs \longleftrightarrow \text{bin-last (bl-to-bin-aux xs w)}$$

⟨proof⟩

lemma *last-bin-last*:

$$\text{size } xs > 0 \implies \text{last } xs \longleftrightarrow \text{bin-last (bl-to-bin xs)}$$

⟨proof⟩

lemma *bin-last-last*:

$$\text{bin-last } w \longleftrightarrow \text{last (bin-to-bl (Suc n) w)}$$

⟨proof⟩

lemma *bl-xor-aux-bin*:

$$\text{map2 } (\%x \ y. x \sim = y) \text{ (bin-to-bl-aux } n \text{ v bs) (bin-to-bl-aux } n \text{ w cs) =}$$

$$\text{bin-to-bl-aux } n \text{ (v XOR w) (map2 } (\%x \ y. x \sim = y) \text{ bs cs)}$$

⟨proof⟩

lemma *bl-or-aux-bin*:

$$\begin{aligned} \text{map2 } (op \mid) (bin\text{-to}\text{-bl}\text{-aux } n \ v \ bs) (bin\text{-to}\text{-bl}\text{-aux } n \ w \ cs) = \\ bin\text{-to}\text{-bl}\text{-aux } n \ (v \ OR \ w) (map2 \ (op \mid) \ bs \ cs) \\ \langle proof \rangle \end{aligned}$$

lemma *bl-and-aux-bin*:

$$\begin{aligned} \text{map2 } (op \ \&) (bin\text{-to}\text{-bl}\text{-aux } n \ v \ bs) (bin\text{-to}\text{-bl}\text{-aux } n \ w \ cs) = \\ bin\text{-to}\text{-bl}\text{-aux } n \ (v \ AND \ w) (map2 \ (op \ \&) \ bs \ cs) \\ \langle proof \rangle \end{aligned}$$

lemma *bl-not-aux-bin*:

$$\begin{aligned} \text{map } Not \ (bin\text{-to}\text{-bl}\text{-aux } n \ w \ cs) = \\ bin\text{-to}\text{-bl}\text{-aux } n \ (NOT \ w) (map \ Not \ cs) \\ \langle proof \rangle \end{aligned}$$

lemma *bl-not-bin*: $\text{map } Not \ (bin\text{-to}\text{-bl } n \ w) = bin\text{-to}\text{-bl } n \ (NOT \ w)$
 $\langle proof \rangle$

lemma *bl-and-bin*:

$$\begin{aligned} \text{map2 } (op \ \wedge) (bin\text{-to}\text{-bl } n \ v) (bin\text{-to}\text{-bl } n \ w) = bin\text{-to}\text{-bl } n \ (v \ AND \ w) \\ \langle proof \rangle \end{aligned}$$

lemma *bl-or-bin*:

$$\begin{aligned} \text{map2 } (op \ \vee) (bin\text{-to}\text{-bl } n \ v) (bin\text{-to}\text{-bl } n \ w) = bin\text{-to}\text{-bl } n \ (v \ OR \ w) \\ \langle proof \rangle \end{aligned}$$

lemma *bl-xor-bin*:

$$\begin{aligned} \text{map2 } (\lambda x \ y. \ x \neq \ y) (bin\text{-to}\text{-bl } n \ v) (bin\text{-to}\text{-bl } n \ w) = bin\text{-to}\text{-bl } n \ (v \ XOR \ w) \\ \langle proof \rangle \end{aligned}$$

lemma *drop-bin2bl-aux*:

$$\begin{aligned} \text{drop } m \ (bin\text{-to}\text{-bl}\text{-aux } n \ bin \ bs) = \\ bin\text{-to}\text{-bl}\text{-aux } (n - m) \ bin \ (drop \ (m - n) \ bs) \\ \langle proof \rangle \end{aligned}$$

lemma *drop-bin2bl*: $\text{drop } m \ (bin\text{-to}\text{-bl } n \ bin) = bin\text{-to}\text{-bl } (n - m) \ bin$
 $\langle proof \rangle$

lemma *take-bin2bl-lem1*:

$$\begin{aligned} \text{take } m \ (bin\text{-to}\text{-bl}\text{-aux } m \ w \ bs) = bin\text{-to}\text{-bl } m \ w \\ \langle proof \rangle \end{aligned}$$

lemma *take-bin2bl-lem*:

$$\begin{aligned} \text{take } m \ (bin\text{-to}\text{-bl}\text{-aux } (m + n) \ w \ bs) = \\ \text{take } m \ (bin\text{-to}\text{-bl } (m + n) \ w) \\ \langle proof \rangle \end{aligned}$$

lemma *bin-split-take*:

$$\begin{aligned} bin\text{-split } n \ c = (a, b) \implies \\ bin\text{-to}\text{-bl } m \ a = \text{take } m \ (bin\text{-to}\text{-bl } (m + n) \ c) \end{aligned}$$

<proof>

lemma *bin-split-take1*:

$k = m + n \implies \text{bin-split } n \ c = (a, b) \implies$
 $\text{bin-to-bl } m \ a = \text{take } m \ (\text{bin-to-bl } k \ c)$
<proof>

lemma *nth-takefill*: $m < n \implies$

$\text{takefill fill } n \ l \ ! \ m = (\text{if } m < \text{length } l \ \text{then } l \ ! \ m \ \text{else } \text{fill})$
<proof>

lemma *takefill-alt*:

$\text{takefill fill } n \ l = \text{take } n \ l \ @ \ \text{replicate } (n - \text{length } l) \ \text{fill}$
<proof>

lemma *takefill-replicate* [*simp*]:

$\text{takefill fill } n \ (\text{replicate } m \ \text{fill}) = \text{replicate } n \ \text{fill}$
<proof>

lemma *takefill-le'*:

$n = m + k \implies \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$
<proof>

lemma *length-takefill* [*simp*]: $\text{length } (\text{takefill fill } n \ l) = n$

<proof>

lemma *take-takefill'*:

$!!w \ n. \ n = k + m \implies \text{take } k \ (\text{takefill fill } n \ w) = \text{takefill fill } k \ w$
<proof>

lemma *drop-takefill*:

$!!w. \ \text{drop } k \ (\text{takefill fill } (m + k) \ w) = \text{takefill fill } m \ (\text{drop } k \ w)$
<proof>

lemma *takefill-le* [*simp*]:

$m \leq n \implies \text{takefill } x \ m \ (\text{takefill } x \ n \ l) = \text{takefill } x \ m \ l$
<proof>

lemma *take-takefill* [*simp*]:

$m \leq n \implies \text{take } m \ (\text{takefill fill } n \ w) = \text{takefill fill } m \ w$
<proof>

lemma *takefill-append*:

$\text{takefill fill } (m + \text{length } xs) \ (xs \ @ \ w) = xs \ @ \ (\text{takefill fill } m \ w)$
<proof>

lemma *takefill-same'*:

$l = \text{length } xs \implies \text{takefill fill } l \ xs = xs$
<proof>

lemmas *takefill-same* [simp] = *takefill-same'* [OF refl]

lemma *takefill-bintrunc*:

takefill False n bl = *rev (bin-to-bl n (bl-to-bin (rev bl)))*
 ⟨proof⟩

lemma *bl-bin-bl-rtf*:

bin-to-bl n (bl-to-bin bl) = *rev (takefill False n (rev bl))*
 ⟨proof⟩

lemma *bl-bin-bl-rep-drop*:

bin-to-bl n (bl-to-bin bl) =
replicate (n - length bl) False @ drop (length bl - n) bl
 ⟨proof⟩

lemma *tf-rev*:

$n + k = m + \text{length } bl \implies \text{takefill } x \ m \ (\text{rev } (\text{takefill } y \ n \ bl)) =$
 $\text{rev } (\text{takefill } y \ m \ (\text{rev } (\text{takefill } x \ k \ (\text{rev } bl))))$
 ⟨proof⟩

lemma *takefill-minus*:

$0 < n \implies \text{takefill fill } (\text{Suc } (n - 1)) \ w = \text{takefill fill } n \ w$
 ⟨proof⟩

lemmas *takefill-Suc-cases* =

list.cases [THEN takefill.Suc [THEN trans]]

lemmas *takefill-Suc-Nil* = *takefill-Suc-cases* (1)

lemmas *takefill-Suc-Cons* = *takefill-Suc-cases* (2)

lemmas *takefill-minus-simps* = *takefill-Suc-cases [THEN [2]*

takefill-minus [symmetric, THEN trans]]

lemma *takefill-numeral-Nil* [simp]:

takefill fill (numeral k) [] = *fill # takefill fill (pred-numeral k) []*
 ⟨proof⟩

lemma *takefill-numeral-Cons* [simp]:

takefill fill (numeral k) (x # xs) = *x # takefill fill (pred-numeral k) xs*
 ⟨proof⟩

lemma *bl-to-bin-aux-cat*:

!!*nv v. bl-to-bin-aux bs (bin-cat w nv v)* =
bin-cat w (nv + length bs) (bl-to-bin-aux bs v)
 ⟨proof⟩

lemma *bin-to-bl-aux-cat*:

$$\begin{aligned} !!w \text{ bs. } \text{bin-to-bl-aux } (nv + nw) (\text{bin-cat } v \text{ nw } w) \text{ bs} = \\ \text{bin-to-bl-aux } nv \text{ v } (\text{bin-to-bl-aux } nw \text{ w } \text{ bs}) \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *bl-to-bin-aux-alt*:

$$\begin{aligned} \text{bl-to-bin-aux } \text{bs } w = \text{bin-cat } w (\text{length } \text{bs}) (\text{bl-to-bin } \text{bs}) \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *bin-to-bl-cat*:

$$\begin{aligned} \text{bin-to-bl } (nv + nw) (\text{bin-cat } v \text{ nw } w) = \\ \text{bin-to-bl-aux } nv \text{ v } (\text{bin-to-bl } nw \text{ w}) \\ \langle \text{proof} \rangle \end{aligned}$$

lemmas *bl-to-bin-aux-app-cat* =

$$\text{trans } [\text{OF } \text{bl-to-bin-aux-append } \text{bl-to-bin-aux-alt}]$$

lemmas *bin-to-bl-aux-cat-app* =

$$\text{trans } [\text{OF } \text{bin-to-bl-aux-cat } \text{bin-to-bl-aux-alt}]$$

lemma *bl-to-bin-app-cat*:

$$\begin{aligned} \text{bl-to-bin } (\text{bsa } @ \text{ bs}) = \text{bin-cat } (\text{bl-to-bin } \text{bsa}) (\text{length } \text{bs}) (\text{bl-to-bin } \text{bs}) \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *bin-to-bl-cat-app*:

$$\begin{aligned} \text{bin-to-bl } (n + nw) (\text{bin-cat } w \text{ nw } wa) = \text{bin-to-bl } n \text{ w } @ \text{bin-to-bl } nw \text{ wa} \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *bl-to-bin-app-cat-alt*:

$$\begin{aligned} \text{bin-cat } (\text{bl-to-bin } \text{cs}) n \text{ w} = \text{bl-to-bin } (\text{cs } @ \text{bin-to-bl } n \text{ w}) \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *mask-lem*: $(\text{bl-to-bin } (\text{True } \# \text{ replicate } n \text{ False})) =$

$$\begin{aligned} (\text{bl-to-bin } (\text{replicate } n \text{ True})) + 1 \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *length-bl-of-nth* [simp]: $\text{length } (\text{bl-of-nth } n \text{ f}) = n$

$$\langle \text{proof} \rangle$$

lemma *nth-bl-of-nth* [simp]:

$$\begin{aligned} m < n \implies \text{rev } (\text{bl-of-nth } n \text{ f}) ! m = f \text{ m} \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *bl-of-nth-inj*:

$$\begin{aligned} (!!k. k < n \implies f \text{ k} = g \text{ k}) \implies \text{bl-of-nth } n \text{ f} = \text{bl-of-nth } n \text{ g} \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *bl-of-nth-nth-le*:

$n \leq \text{length } xs \implies \text{bl-of-nth } n \ (\text{nth } (\text{rev } xs)) = \text{drop } (\text{length } xs - n) \ xs$
 ⟨proof⟩

lemma *bl-of-nth-nth [simp]*: $\text{bl-of-nth } (\text{length } xs) \ (\text{op } ! \ (\text{rev } xs)) = xs$

⟨proof⟩

lemma *size-rbl-pred*: $\text{length } (\text{rbl-pred } bl) = \text{length } bl$

⟨proof⟩

lemma *size-rbl-succ*: $\text{length } (\text{rbl-succ } bl) = \text{length } bl$

⟨proof⟩

lemma *size-rbl-add*:

!!cl. $\text{length } (\text{rbl-add } bl \ cl) = \text{length } bl$

⟨proof⟩

lemma *size-rbl-mult*:

!!cl. $\text{length } (\text{rbl-mult } bl \ cl) = \text{length } bl$

⟨proof⟩

lemmas *rbl-sizes [simp]* =

size-rbl-pred size-rbl-succ size-rbl-add size-rbl-mult

lemmas *rbl-Nils* =

rbl-pred.Nil rbl-succ.Nil rbl-add.Nil rbl-mult.Nil

lemma *rbl-pred*:

$\text{rbl-pred } (\text{rev } (\text{bin-to-bl } n \ bin)) = \text{rev } (\text{bin-to-bl } n \ (bin - 1))$

⟨proof⟩

lemma *rbl-succ*:

$\text{rbl-succ } (\text{rev } (\text{bin-to-bl } n \ bin)) = \text{rev } (\text{bin-to-bl } n \ (bin + 1))$

⟨proof⟩

lemma *rbl-add*:

!!bina binb. $\text{rbl-add } (\text{rev } (\text{bin-to-bl } n \ bina)) \ (\text{rev } (\text{bin-to-bl } n \ binb)) =$
 $\text{rev } (\text{bin-to-bl } n \ (bina + binb))$

⟨proof⟩

lemma *rbl-add-app2*:

!!blb. $\text{length } blb \geq \text{length } bla \implies$

$\text{rbl-add } bla \ (blb @ blc) = \text{rbl-add } bla \ blb$

⟨proof⟩

lemma *rbl-add-take2*:

!!blb. $\text{length } blb \geq \text{length } bla \implies$

$\text{rbl-add } bla \ (\text{take } (\text{length } bla) \ blb) = \text{rbl-add } bla \ blb$

⟨proof⟩

lemma *rbl-add-long*:

$$m \geq n \implies \text{rbl-add } (\text{rev } (\text{bin-to-bl } n \text{ bina})) (\text{rev } (\text{bin-to-bl } m \text{ binb})) = \\ \text{rev } (\text{bin-to-bl } n \text{ (bina + binb)}) \\ \langle \text{proof} \rangle$$

lemma *rbl-mult-app2*:

$$!!\text{blb. length blb} \geq \text{length bla} \implies \\ \text{rbl-mult bla (blb @ blc)} = \text{rbl-mult bla blb} \\ \langle \text{proof} \rangle$$

lemma *rbl-mult-take2*:

$$\text{length blb} \geq \text{length bla} \implies \\ \text{rbl-mult bla (take (length bla) blb)} = \text{rbl-mult bla blb} \\ \langle \text{proof} \rangle$$

lemma *rbl-mult-gt1*:

$$m \geq \text{length bl} \implies \text{rbl-mult bl (rev (bin-to-bl } m \text{ binb}))} = \\ \text{rbl-mult bl (rev (bin-to-bl (length bl) binb))} \\ \langle \text{proof} \rangle$$

lemma *rbl-mult-gt*:

$$m > n \implies \text{rbl-mult } (\text{rev } (\text{bin-to-bl } n \text{ bina})) (\text{rev } (\text{bin-to-bl } m \text{ binb})) = \\ \text{rbl-mult } (\text{rev } (\text{bin-to-bl } n \text{ bina})) (\text{rev } (\text{bin-to-bl } n \text{ binb})) \\ \langle \text{proof} \rangle$$

lemmas *rbl-mult-Suc = lessI [THEN rbl-mult-gt]*

lemma *rbbL-Cons*:

$$b \# \text{rev } (\text{bin-to-bl } n \text{ x}) = \text{rev } (\text{bin-to-bl } (\text{Suc } n) \text{ (x BIT b)}) \\ \langle \text{proof} \rangle$$

lemma *rbl-mult: !!bina binb.*

$$\text{rbl-mult } (\text{rev } (\text{bin-to-bl } n \text{ bina})) (\text{rev } (\text{bin-to-bl } n \text{ binb})) = \\ \text{rev } (\text{bin-to-bl } n \text{ (bina * binb)}) \\ \langle \text{proof} \rangle$$

lemma *rbl-add-split*:

$$P (\text{rbl-add } (y \# ys) (x \# xs)) = \\ (\text{ALL } ws. \text{length } ws = \text{length } ys \implies ws = \text{rbl-add } ys \text{ xs} \implies \\ (y \implies ((x \implies P (\text{False} \# \text{rbl-succ } ws)) \ \& \ (\sim x \implies P (\text{True} \# ws)))) \ \& \\ (\sim y \implies P (x \# ws))) \\ \langle \text{proof} \rangle$$

lemma *rbl-mult-split*:

$$P (\text{rbl-mult } (y \# ys) \text{ xs}) = \\ (\text{ALL } ws. \text{length } ws = \text{Suc } (\text{length } ys) \implies ws = \text{False} \# \text{rbl-mult } ys \text{ xs} \implies \\ (y \implies P (\text{rbl-add } ws \text{ xs})) \ \& \ (\sim y \implies P \text{ ws}))$$

$\langle \text{proof} \rangle$

12.3 Repeated splitting or concatenation

lemma *sclm*:

$\text{size} (\text{concat} (\text{map} (\text{bin-to-bl } n) xs)) = \text{length } xs * n$
 $\langle \text{proof} \rangle$

lemma *bin-cat-foldl-lem*:

$\text{foldl} (\%u. \text{bin-cat } u \ n) \ x \ xs =$
 $\text{bin-cat } x \ (\text{size } xs * n) \ (\text{foldl} (\%u. \text{bin-cat } u \ n) \ y \ xs)$
 $\langle \text{proof} \rangle$

lemma *bin-rcat-bl*:

$(\text{bin-rcat } n \ wl) = \text{bl-to-bin} (\text{concat} (\text{map} (\text{bin-to-bl } n) wl))$
 $\langle \text{proof} \rangle$

lemmas *bin-rsplit-aux-simps* = *bin-rsplit-aux.simps* *bin-rsplittl-aux.simps*

lemmas *rsplit-aux-simps* = *bin-rsplit-aux-simps*

lemmas *th-if-simp1* = *if-split* [where $P = op = l$, THEN *iffD1*, THEN *conjunct1*, THEN *mp*] for l

lemmas *th-if-simp2* = *if-split* [where $P = op = l$, THEN *iffD1*, THEN *conjunct2*, THEN *mp*] for l

lemmas *rsplit-aux-simp1s* = *rsplit-aux-simps* [THEN *th-if-simp1*]

lemmas *rsplit-aux-simp2ls* = *rsplit-aux-simps* [THEN *th-if-simp2*]

lemmas *bin-rsplit-aux-simp2s* [*simp*] = *rsplit-aux-simp2ls* [*unfolded Let-def*]

lemmas *rbscl* = *bin-rsplit-aux-simp2s* (2)

lemmas *rsplit-aux-0-simps* [*simp*] =

rsplit-aux-simp1s [*OF disjI1*] *rsplit-aux-simp1s* [*OF disjI2*]

lemma *bin-rsplit-aux-append*:

$\text{bin-rsplit-aux } n \ m \ c \ (bs \ @ \ cs) = \text{bin-rsplit-aux } n \ m \ c \ bs \ @ \ cs$
 $\langle \text{proof} \rangle$

lemma *bin-rsplittl-aux-append*:

$\text{bin-rsplittl-aux } n \ m \ c \ (bs \ @ \ cs) = \text{bin-rsplittl-aux } n \ m \ c \ bs \ @ \ cs$
 $\langle \text{proof} \rangle$

lemmas *rsplit-aux-apps* [where $bs = []$] =

bin-rsplit-aux-append *bin-rsplittl-aux-append*

lemmas *rsplit-def-auxs* = *bin-rsplit-def* *bin-rsplittl-def*

lemmas *rsplit-aux-alts* = *rsplit-aux-apps*

[*unfolded append-Nil rsplit-def-auxs [symmetric]*]

lemma *bin-split-minus*: $0 < n \implies \text{bin-split } (\text{Suc } (n - 1)) w = \text{bin-split } n w$
 ⟨*proof*⟩

lemmas *bin-split-minus-simp* =
bin-split.Suc [THEN [2] bin-split-minus [symmetric, THEN trans]]

lemma *bin-split-pred-simp* [*simp*]:
 $(0::\text{nat}) < \text{numeral } \text{bin} \implies$
 $\text{bin-split } (\text{numeral } \text{bin}) w =$
 $(\text{let } (w1, w2) = \text{bin-split } (\text{numeral } \text{bin} - 1) (\text{bin-rest } w)$
 $\text{in } (w1, w2 \text{ BIT } \text{bin-last } w))$
 ⟨*proof*⟩

lemma *bin-rsplit-aux-simp-alt*:
 $\text{bin-rsplit-aux } n m c \text{ bs} =$
 $(\text{if } m = 0 \vee n = 0$
 $\text{then } \text{bs}$
 $\text{else let } (a, b) = \text{bin-split } n c \text{ in } \text{bin-rsplit } n (m - n, a) @ b \# \text{bs})$
 ⟨*proof*⟩

lemmas *bin-rsplit-simp-alt* =
trans [OF bin-rsplit-def bin-rsplit-aux-simp-alt]

lemmas *bthrs* = *bin-rsplit-simp-alt [THEN [2] trans]*

lemma *bin-rsplit-size-sign'* [*rule-format*] :
 $[[n > 0; \text{rev } sw = \text{bin-rsplit } n (nw, w)] \implies$
 $(\text{ALL } v: \text{set } sw. \text{bintrunc } n v = v)$
 ⟨*proof*⟩

lemmas *bin-rsplit-size-sign* = *bin-rsplit-size-sign' [OF asm-rl*
rev-rev-ident [THEN trans] set-rev [THEN equalityD2 [THEN subsetD]]]

lemma *bin-nth-rsplit* [*rule-format*] :
 $n > 0 \implies m < n \implies (\text{ALL } w k nw. \text{rev } sw = \text{bin-rsplit } n (nw, w) \dashrightarrow$
 $k < \text{size } sw \dashrightarrow \text{bin-nth } (sw ! k) m = \text{bin-nth } w (k * n + m))$
 ⟨*proof*⟩

lemma *bin-rsplit-all*:
 $0 < nw \implies nw \leq n \implies \text{bin-rsplit } n (nw, w) = [\text{bintrunc } n w]$
 ⟨*proof*⟩

lemma *bin-rsplit-l* [*rule-format*] :
 $\text{ALL } \text{bin}. \text{bin-rsplitl } n (m, \text{bin}) = \text{bin-rsplit } n (m, \text{bintrunc } m \text{ bin})$
 ⟨*proof*⟩

lemma *bin-rsplit-rcat* [*rule-format*] :

$n > 0 \dashrightarrow \text{bin-rsplit } n (n * \text{size } ws, \text{bin-rcat } n \text{ } ws) = \text{map } (\text{bintrunc } n) \text{ } ws$
 ⟨proof⟩

lemma *bin-rsplit-aux-len-le* [rule-format] :
 $\forall ws \ m. \ n \neq 0 \longrightarrow ws = \text{bin-rsplit-aux } n \ nw \ w \ bs \longrightarrow$
 $\text{length } ws \leq m \longleftrightarrow nw + \text{length } bs * n \leq m * n$
 ⟨proof⟩

lemma *bin-rsplit-len-le*:
 $n \neq 0 \dashrightarrow ws = \text{bin-rsplit } n (nw, w) \dashrightarrow (\text{length } ws \leq m) = (nw \leq m * n)$
 ⟨proof⟩

lemma *bin-rsplit-aux-len*:
 $n \neq 0 \implies \text{length } (\text{bin-rsplit-aux } n \ nw \ w \ cs) =$
 $(nw + n - 1) \text{ div } n + \text{length } cs$
 ⟨proof⟩

lemma *bin-rsplit-len*:
 $n \neq 0 \implies \text{length } (\text{bin-rsplit } n (nw, w)) = (nw + n - 1) \text{ div } n$
 ⟨proof⟩

lemma *bin-rsplit-aux-len-indep*:
 $n \neq 0 \implies \text{length } bs = \text{length } cs \implies$
 $\text{length } (\text{bin-rsplit-aux } n \ nw \ v \ bs) =$
 $\text{length } (\text{bin-rsplit-aux } n \ nw \ w \ cs)$
 ⟨proof⟩

lemma *bin-rsplit-len-indep*:
 $n \neq 0 \implies \text{length } (\text{bin-rsplit } n (nw, v)) = \text{length } (\text{bin-rsplit } n (nw, w))$
 ⟨proof⟩

Even more bit operations

instantiation *int* :: *bitss*
begin

definition [*iff*]:
 $i \ !! \ n \longleftrightarrow \text{bin-nth } i \ n$

definition
 $\text{lsb } i = (i \ !! \ 0)$

definition
 $\text{set-bit } i \ n \ b = \text{bin-sc } n \ b \ i$

definition
 $\text{set-bits } f =$
 (if $\exists n. \forall n' \geq n. \neg f \ n'$ then
 let $n = \text{LEAST } n. \forall n' \geq n. \neg f \ n'$

```

    in bl-to-bin (rev (map f [0.. $n$ ]))
  else if  $\exists n. \forall n' \geq n. f n'$  then
    let  $n = \text{LEAST } n. \forall n' \geq n. f n'$ 
    in sbintrunc  $n$  (bl-to-bin (True # rev (map f [0.. $n$ ])))
  else 0 :: int)

```

definition

$$\text{shiffl } x \ n = (x :: \text{int}) * 2 \wedge n$$
definition

$$\text{shiftr } x \ n = (x :: \text{int}) \text{ div } 2 \wedge n$$
definition

$$\text{msb } x \longleftrightarrow (x :: \text{int}) < 0$$
instance $\langle \text{proof} \rangle$ **end****end**

13 Type Definition Theorems

theory *Misc-Typedef***imports** *Main***begin**

14 More lemmas about normal type definitions

lemma

```

tdD1: type-definition Rep Abs A  $\implies \forall x. \text{Rep } x \in A$  and
tdD2: type-definition Rep Abs A  $\implies \forall x. \text{Abs } (\text{Rep } x) = x$  and
tdD3: type-definition Rep Abs A  $\implies \forall y. y \in A \longrightarrow \text{Rep } (\text{Abs } y) = y$ 
 $\langle \text{proof} \rangle$ 

```

lemma *td-nat-int*:

```

type-definition int nat (Collect (op <= 0))
 $\langle \text{proof} \rangle$ 

```

context *type-definition***begin****declare** *Rep* [*iff*] *Rep-inverse* [*simp*] *Rep-inject* [*simp*]

```

lemma Abs-eqD:  $\text{Abs } x = \text{Abs } y \implies x \in A \implies y \in A \implies x = y$ 
 $\langle \text{proof} \rangle$ 

```

lemma *Abs-inverse'*:

$r : A \implies \text{Abs } r = a \implies \text{Rep } a = r$
 ⟨proof⟩

lemma *Rep-comp-inverse*:

$\text{Rep } o f = g \implies \text{Abs } o g = f$
 ⟨proof⟩

lemma *Rep-eqD* [elim!]: $\text{Rep } x = \text{Rep } y \implies x = y$
 ⟨proof⟩

lemma *Rep-inverse'*: $\text{Rep } a = r \implies \text{Abs } r = a$
 ⟨proof⟩

lemma *comp-Abs-inverse*:

$f o \text{Abs} = g \implies g o \text{Rep} = f$
 ⟨proof⟩

lemma *set-Rep*:

$A = \text{range } \text{Rep}$
 ⟨proof⟩

lemma *set-Rep-Abs*: $A = \text{range } (\text{Rep } o \text{Abs})$
 ⟨proof⟩

lemma *Abs-inj-on*: *inj-on* $\text{Abs } A$
 ⟨proof⟩

lemma *image*: $\text{Abs } ` A = \text{UNIV}$
 ⟨proof⟩

lemmas *td-thm* = *type-definition-axioms*

lemma *fns1*:

$\text{Rep } o fa = fr o \text{Rep} \mid fa o \text{Abs} = \text{Abs } o fr \implies \text{Abs } o fr o \text{Rep} = fa$
 ⟨proof⟩

lemmas *fns1a* = *disjI1* [THEN *fns1*]

lemmas *fns1b* = *disjI2* [THEN *fns1*]

lemma *fns4*:

$\text{Rep } o fa o \text{Abs} = fr \implies$
 $\text{Rep } o fa = fr o \text{Rep} \ \& \ fa o \text{Abs} = \text{Abs } o fr$
 ⟨proof⟩

end

interpretation *nat-int*: *type-definition int nat Collect* (*op* <= 0)
 ⟨proof⟩

declare

nat-int.Rep-cases [*cases del*]
nat-int.Abs-cases [*cases del*]
nat-int.Rep-induct [*induct del*]
nat-int.Abs-induct [*induct del*]

14.1 Extended form of type definition predicate

lemma *td-conds*:

$norm \circ norm = norm \implies (fr \circ norm = norm \circ fr) =$
 $(norm \circ fr \circ norm = fr \circ norm \ \& \ norm \circ fr \circ norm = norm \circ fr)$
 $\langle proof \rangle$

lemma *fn-comm-power*:

$fa \circ tr = tr \circ fr \implies fa \wedge n \circ tr = tr \circ fr \wedge n$
 $\langle proof \rangle$

lemmas *fn-comm-power'* =

ext [THEN *fn-comm-power*, THEN *fun-cong*, *unfolded o-def*]

locale *td-ext* = *type-definition* +

fixes *norm*

assumes *eq-norm*: $\bigwedge x. Rep (Abs x) = norm x$

begin

lemma *Abs-norm* [*simp*]:

$Abs (norm x) = Abs x$
 $\langle proof \rangle$

lemma *td-th*:

$g \circ Abs = f \implies f (Rep x) = g x$
 $\langle proof \rangle$

lemma *eq-norm'*: $Rep \circ Abs = norm$

$\langle proof \rangle$

lemma *norm-Rep* [*simp*]: $norm (Rep x) = Rep x$

$\langle proof \rangle$

lemmas *td* = *td-thm*

lemma *set-iff-norm*: $w : A \longleftrightarrow w = norm w$

$\langle proof \rangle$

lemma *inverse-norm*:

$(Abs n = w) = (Rep w = norm n)$
 $\langle proof \rangle$

lemma *norm-eq-iff*:

$(\text{norm } x = \text{norm } y) = (\text{Abs } x = \text{Abs } y)$
 $\langle \text{proof} \rangle$

lemma *norm-comps*:

$\text{Abs } o \text{ norm} = \text{Abs}$
 $\text{norm } o \text{ Rep} = \text{Rep}$
 $\text{norm } o \text{ norm} = \text{norm}$
 $\langle \text{proof} \rangle$

lemmas *norm-norm* [*simp*] = *norm-comps*

lemma *fn5*:

$\text{Rep } o \text{ fa } o \text{ Abs} = \text{fr} ==>$
 $\text{fr } o \text{ norm} = \text{fr} \ \& \ \text{norm } o \text{ fr} = \text{fr}$
 $\langle \text{proof} \rangle$

lemma *fn2*:

$\text{Abs } o \text{ fr } o \text{ Rep} = \text{fa} ==>$
 $(\text{norm } o \text{ fr } o \text{ norm} = \text{fr } o \text{ norm}) = (\text{Rep } o \text{ fa} = \text{fr } o \text{ Rep})$
 $\langle \text{proof} \rangle$

lemma *fn3*:

$\text{Abs } o \text{ fr } o \text{ Rep} = \text{fa} ==>$
 $(\text{norm } o \text{ fr } o \text{ norm} = \text{norm } o \text{ fr}) = (\text{fa } o \text{ Abs} = \text{Abs } o \text{ fr})$
 $\langle \text{proof} \rangle$

lemma *fn4*:

$\text{fr } o \text{ norm} = \text{norm } o \text{ fr} ==>$
 $(\text{fa } o \text{ Abs} = \text{Abs } o \text{ fr}) = (\text{Rep } o \text{ fa} = \text{fr } o \text{ Rep})$
 $\langle \text{proof} \rangle$

lemma *range-norm*:

$\text{range } (\text{Rep } o \text{ Abs}) = A$
 $\langle \text{proof} \rangle$

end

lemmas *td-ext-def'* =

td-ext-def [*unfolded type-definition-def* *td-ext-axioms-def*]

end

15 Miscellaneous lemmas, of at least doubtful value

theory *Word-Miscellaneous*

imports *Main* $\sim\sim$ */src/HOL/Library/Bit Misc-Numeric*

begin

lemma *power-minus-simp*:

$$0 < n \implies a ^ n = a * a ^ (n - 1)$$

<proof>

lemma *funpow-minus-simp*:

$$0 < n \implies f ^ n = f \circ f ^ (n - 1)$$

<proof>

lemma *power-numeral*:

$$a ^ numeral\ k = a * a ^ (pred\ numeral\ k)$$

<proof>

lemma *funpow-numeral [simp]*:

$$f ^ numeral\ k = f \circ f ^ (pred\ numeral\ k)$$

<proof>

lemma *replicate-numeral [simp]*:

$$replicate\ (numeral\ k)\ x = x \# replicate\ (pred\ numeral\ k)\ x$$

<proof>

lemma *rco-alt*: $(f \circ g) ^ n \circ f = f \circ (g \circ f) ^ n$

<proof>

lemma *list-exhaust-size-gt0*:

assumes $y: \bigwedge a\ list. y = a \# list \implies P$
shows $0 < length\ y \implies P$
<proof>

lemma *list-exhaust-size-eq0*:

assumes $y: y = [] \implies P$
shows $length\ y = 0 \implies P$
<proof>

lemma *size-Cons-lem-eq*:

$$y = xa \# list \implies size\ y = Suc\ k \implies size\ list = k$$

<proof>

lemmas *ls-splits = prod.split prod.split-asm if-split-asm*

lemma *not-B1-is-B0*: $y \neq (1::bit) \implies y = (0::bit)$

<proof>

lemma *B1-ass-B0*:

assumes $y: y = (0::bit) \implies y = (1::bit)$
shows $y = (1::bit)$
<proof>

lemmas *n2s-ths [THEN eq-reflection] = add-2-eq-Suc add-2-eq-Suc'*

lemmas $s2n\text{-ths} = n2s\text{-ths}$ [symmetric]

lemma *and-len*: $xs = ys \implies xs = ys \ \& \ \text{length } xs = \text{length } ys$
 ⟨proof⟩

lemma *size-if*: $\text{size } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{size } xs \text{ else } \text{size } ys)$
 ⟨proof⟩

lemma *tl-if*: $\text{tl } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{tl } xs \text{ else } \text{tl } ys)$
 ⟨proof⟩

lemma *hd-if*: $\text{hd } (\text{if } p \text{ then } xs \text{ else } ys) = (\text{if } p \text{ then } \text{hd } xs \text{ else } \text{hd } ys)$
 ⟨proof⟩

lemma *if-Not-x*: $(\text{if } p \text{ then } \sim x \text{ else } x) = (p = (\sim x))$
 ⟨proof⟩

lemma *if-x-Not*: $(\text{if } p \text{ then } x \text{ else } \sim x) = (p = x)$
 ⟨proof⟩

lemma *if-same-and*: $(\text{If } p \ x \ y \ \& \ \text{If } p \ u \ v) = (\text{if } p \ \text{then } x \ \& \ u \ \text{else } y \ \& \ v)$
 ⟨proof⟩

lemma *if-same-eq*: $(\text{If } p \ x \ y = \text{If } p \ u \ v) = (\text{if } p \ \text{then } x = (u) \ \text{else } y = (v))$
 ⟨proof⟩

lemma *if-same-eq-not*:
 $(\text{If } p \ x \ y = (\sim \text{If } p \ u \ v)) = (\text{if } p \ \text{then } x = (\sim u) \ \text{else } y = (\sim v))$
 ⟨proof⟩

lemma *if-Cons*: $(\text{if } p \ \text{then } x \ \# \ xs \ \text{else } y \ \# \ ys) = \text{If } p \ x \ y \ \# \ \text{If } p \ xs \ ys$
 ⟨proof⟩

lemma *if-single*:
 $(\text{if } xc \ \text{then } [xab] \ \text{else } [an]) = [\text{if } xc \ \text{then } xab \ \text{else } an]$
 ⟨proof⟩

lemma *if-bool-simps*:
 $\text{If } p \ \text{True } y = (p \ | \ y) \ \& \ \text{If } p \ \text{False } y = (\sim p \ \& \ y) \ \&$
 $\text{If } p \ y \ \text{True} = (p \ \dashrightarrow y) \ \& \ \text{If } p \ y \ \text{False} = (p \ \& \ y)$
 ⟨proof⟩

lemmas *if-simps* = *if-x-Not if-Not-x if-cancel if-True if-False if-bool-simps*

lemmas *seqr* = *eq-reflection* [where $x = \text{size } w$] for w

lemma *the-elmI*: $y = \{x\} \implies \text{the-elm } y = x$
 ⟨proof⟩

lemma *nonemptyE*: $S \sim = \{\}$ $\implies (\exists x. x : S \implies R) \implies R$ $\langle proof \rangle$

lemma *gt-or-eq-0*: $0 < y \vee 0 = (y :: nat)$ $\langle proof \rangle$

lemmas *xtr1* = *xtrans*(1)

lemmas *xtr2* = *xtrans*(2)

lemmas *xtr3* = *xtrans*(3)

lemmas *xtr4* = *xtrans*(4)

lemmas *xtr5* = *xtrans*(5)

lemmas *xtr6* = *xtrans*(6)

lemmas *xtr7* = *xtrans*(7)

lemmas *xtr8* = *xtrans*(8)

lemmas *nat-simps* = *diff-add-inverse2* *diff-add-inverse*

lemmas *nat-iffs* = *le-add1* *le-add2*

lemma *sum-imp-diff*: $j = k + i \implies j - i = (k :: nat)$ $\langle proof \rangle$

lemmas *pos-mod-sign2* = *zless2* [*THEN* *pos-mod-sign* [**where** $b = 2 :: int$]]

lemmas *pos-mod-bound2* = *zless2* [*THEN* *pos-mod-bound* [**where** $b = 2 :: int$]]

lemma *nmod2*: $n \bmod (2 :: int) = 0 \mid n \bmod 2 = 1$
 $\langle proof \rangle$

lemmas *eme1p* = *emep1* [*simplified add commute*]

lemma *le-diff-eq'*: $(a \leq c - b) = (b + a \leq (c :: int))$ $\langle proof \rangle$

lemma *less-diff-eq'*: $(a < c - b) = (b + a < (c :: int))$ $\langle proof \rangle$

lemma *diff-less-eq'*: $(a - b < c) = (a < b + (c :: int))$ $\langle proof \rangle$

lemmas *m1mod2k* = *mult-pos-pos* [*OF* *zless2* *zless2p*, *THEN* *zmod-minus1*]

lemma *z1pdiv2*:

$(2 * b + 1) \bmod 2 = (b :: int)$ $\langle proof \rangle$

lemmas *zdiv-le-dividend* = *xtr3* [*OF* *div-by-1* [*symmetric*] *zdiv-mono2*,
simplified int-one-le-iff-zero-less, *simplified*]

lemma *axbyy*:

$a + m + m = b + n + n \implies (a = 0 \mid a = 1) \implies (b = 0 \mid b = 1) \implies$
 $a = b \ \& \ m = (n :: int)$ $\langle proof \rangle$

lemma *axrmod2*:

$(1 + x + x) \bmod 2 = (1 :: int) \ \& \ (0 + x + x) \bmod 2 = (0 :: int)$ $\langle proof \rangle$

lemma *axdiv2*:

$(1 + x + x) \text{ div } 2 = (x :: \text{int}) \ \& \ (0 + x + x) \text{ div } 2 = (x :: \text{int}) \ \langle \text{proof} \rangle$

lemmas *iszero-minus* = *trans* [*THEN trans*,
OF iszero-def neg-equal-0-iff-equal iszero-def [symmetric]]

lemmas *zadd-diff-inverse* = *trans* [*OF diff-add-cancel [symmetric] add.commute*]

lemmas *add-diff-cancel2* = *add.commute* [*THEN diff-eq-eq [THEN iffD2]*]

lemmas *rdmods [symmetric]* = *mod-minus-eq*
mod-diff-left-eq mod-diff-right-eq mod-add-left-eq
mod-add-right-eq mod-mult-right-eq mod-mult-left-eq

lemma *mod-plus-right*:
 $((a + x) \text{ mod } m = (b + x) \text{ mod } m) = (a \text{ mod } m = b \text{ mod } (m :: \text{nat}))$
 $\langle \text{proof} \rangle$

lemma *nat-minus-mod*: $(n - n \text{ mod } m) \text{ mod } m = (0 :: \text{nat})$
 $\langle \text{proof} \rangle$

lemmas *nat-minus-mod-plus-right* = *trans* [*OF nat-minus-mod mod-0 [symmetric]*,
THEN mod-plus-right [THEN iffD2], simplified]

lemmas *push-mods'* = *mod-add-eq*
mod-mult-eq mod-diff-eq
mod-minus-eq

lemmas *push-mods* = *push-mods'* [*THEN eq-reflection*]
lemmas *pull-mods* = *push-mods* [*symmetric*] *rdmods* [*THEN eq-reflection*]
lemmas *mod-simps* =
mod-mult-self2-is-0 [THEN eq-reflection]
mod-mult-self1-is-0 [THEN eq-reflection]
mod-mod-trivial [THEN eq-reflection]

lemma *nat-mod-eq*:
 $!!b. b < n ==> a \text{ mod } n = b \text{ mod } n ==> a \text{ mod } n = (b :: \text{nat})$
 $\langle \text{proof} \rangle$

lemmas *nat-mod-eq'* = *refl* [*THEN [2] nat-mod-eq*]

lemma *nat-mod-lem*:
 $(0 :: \text{nat}) < n ==> b < n = (b \text{ mod } n = b)$
 $\langle \text{proof} \rangle$

lemma *mod-nat-add*:
 $(x :: \text{nat}) < z ==> y < z ==>$
 $(x + y) \text{ mod } z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
 $\langle \text{proof} \rangle$

lemma *mod-nat-sub*:

$(x :: \text{nat}) < z \implies (x - y) \text{ mod } z = x - y$
 ⟨proof⟩

lemma *int-mod-eq*:

$(0 :: \text{int}) \leq b \implies b < n \implies a \text{ mod } n = b \text{ mod } n \implies a \text{ mod } n = b$
 ⟨proof⟩

lemmas *int-mod-eq' = mod-pos-pos-trivial*

lemma *int-mod-le*: $(0 :: \text{int}) \leq a \implies a \text{ mod } n \leq a$

⟨proof⟩

lemma *mod-add-if-z*:

$(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x + y) \text{ mod } z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
 ⟨proof⟩

lemma *mod-sub-if-z*:

$(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x - y) \text{ mod } z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$
 ⟨proof⟩

lemmas *zmde = zmod-zdiv-equality [THEN diff-eq-eq [THEN iffD2], symmetric]*

lemmas *mcl = mult-cancel-left [THEN iffD1, THEN make-pos-rule]*

lemma *zdiv-mult-self*: $m \sim = (0 :: \text{int}) \implies (a + m * n) \text{ div } m = a \text{ div } m + n$

⟨proof⟩

lemma *mod-power-lem*:

$a > 1 \implies a ^ n \text{ mod } a ^ m = (\text{if } m \leq n \text{ then } 0 \text{ else } (a :: \text{int}) ^ n)$
 ⟨proof⟩

lemma *pl-pl-rels*:

$a + b = c + d \implies$
 $a \geq c \ \& \ b \leq d \mid a \leq c \ \& \ b \geq d \implies (d :: \text{nat})$ ⟨proof⟩

lemmas *pl-pl-rels' = add commute [THEN [2] trans, THEN pl-pl-rels]*

lemma *minus-eq*: $(m - k = m) = (k = 0 \mid m = (0 :: \text{nat}))$ ⟨proof⟩

lemma *pl-pl-mm*: $(a :: \text{nat}) + b = c + d \implies a - c = d - b$ ⟨proof⟩

lemmas *pl-pl-mm' = add commute [THEN [2] trans, THEN pl-pl-mm]*

lemmas *dme = box-equals [OF div-mod-equality add-0-right add-0-right]*

lemmas *dtle = xtr3 [OF dme [symmetric] le-add1]*

lemmas *th2 = order-trans [OF order-refl [THEN [2] mult-le-mono] dtle]*

lemma *td-gal*:

$$0 < c \implies (a \geq b * c) = (a \operatorname{div} c \geq (b :: \operatorname{nat}))$$

<proof>

lemmas *td-gal-lt* = *td-gal* [*simplified not-less* [*symmetric*], *simplified*]

lemma *div-mult-le*: $(a :: \operatorname{nat}) \operatorname{div} b * b \leq a$

<proof>

lemmas *sdl* = *split-div-lemma* [*THEN iffD1*, *symmetric*]

lemma *given-quot*: $f > (0 :: \operatorname{nat}) \implies (f * l + (f - 1)) \operatorname{div} f = l$

<proof>

lemma *given-quot-alt*: $f > (0 :: \operatorname{nat}) \implies (l * f + f - \operatorname{Suc} 0) \operatorname{div} f = l$

<proof>

lemma *diff-mod-le*: $(a :: \operatorname{nat}) < d \implies b \operatorname{dvd} d \implies a - a \operatorname{mod} b \leq d - b$

<proof>

lemma *less-le-mult'*:

$$w * c < b * c \implies 0 \leq c \implies (w + 1) * c \leq b * (c :: \operatorname{int})$$

<proof>

lemma *less-le-mult*:

$$w * c < b * c \implies 0 \leq c \implies w * c + c \leq b * (c :: \operatorname{int})$$

<proof>

lemmas *less-le-mult-minus* = *iffD2* [*OF le-diff-eq less-le-mult*, *simplified left-diff-distrib*]

lemma *gen-minus*: $0 < n \implies f n = f (\operatorname{Suc} (n - 1))$

<proof>

lemma *mpl-lem*: $j \leq (i :: \operatorname{nat}) \implies k < j \implies i - j + k < i$ *<proof>*

lemma *nonneg-mod-div*:

$$0 \leq a \implies 0 \leq b \implies 0 \leq (a \operatorname{mod} b :: \operatorname{int}) \ \& \ 0 \leq a \operatorname{div} b$$

<proof>

declare *iszero-0* [*intro*]

lemma *min-pm* [*simp*]:

$$\min a b + (a - b) = (a :: \operatorname{nat})$$

<proof>

lemma *min-pm1* [*simp*]:

$$a - b + \min a b = (a :: \operatorname{nat})$$

<proof>

lemma *rev-min-pm* [*simp*]:
 $\text{min } b \ a + (a - b) = (a :: \text{nat})$
<proof>

lemma *rev-min-pm1* [*simp*]:
 $a - b + \text{min } b \ a = (a :: \text{nat})$
<proof>

lemma *min-minus* [*simp*]:
 $\text{min } m \ (m - k) = (m - k :: \text{nat})$
<proof>

lemma *min-minus'* [*simp*]:
 $\text{min } (m - k) \ m = (m - k :: \text{nat})$
<proof>

end

16 A type of finite bit strings

theory *Word*

imports

Type-Length

~/src/HOL/Library/Boolean-Algebra

Bits-Bit

Bool-List-Representation

Misc-Typedef

Word-Miscellaneous

begin

See *Examples/WordExamples.thy* for examples.

16.1 Type definition

typedef (**overloaded**) *'a word* = $\{(0::\text{int}) ..< 2^{\text{len-of TYPE('a::len0)}\}$
morphisms *uint Abs-word* *<proof>*

lemma *uint-nonnegative*:

$0 \leq \text{uint } w$

<proof>

lemma *uint-bounded*:

fixes *w :: 'a::len0 word*

shows $\text{uint } w < 2^{\text{len-of TYPE('a)}}$

<proof>

lemma *uint-idem*:

fixes $w :: 'a::len0\ word$
shows $uint\ w\ mod\ 2\ ^\ len-of\ TYPE('a) = uint\ w$
 $\langle proof \rangle$

lemma $word-uint-eq-iff$:
 $a = b \longleftrightarrow uint\ a = uint\ b$
 $\langle proof \rangle$

lemma $word-uint-eqI$:
 $uint\ a = uint\ b \implies a = b$
 $\langle proof \rangle$

definition $word-of-int :: int \Rightarrow 'a::len0\ word$
where

— representation of words using unsigned or signed bins, only difference in these is the type class

$word-of-int\ k = Abs-word\ (k\ mod\ 2\ ^\ len-of\ TYPE('a))$

lemma $uint-word-of-int$:
 $uint\ (word-of-int\ k :: 'a::len0\ word) = k\ mod\ 2\ ^\ len-of\ TYPE('a)$
 $\langle proof \rangle$

lemma $word-of-int-uint$:
 $word-of-int\ (uint\ w) = w$
 $\langle proof \rangle$

lemma $split-word-all$:
 $(\bigwedge x::'a::len0\ word. PROP\ P\ x) \equiv (\bigwedge x. PROP\ P\ (word-of-int\ x))$
 $\langle proof \rangle$

16.2 Type conversions and casting

definition $sint :: 'a::len\ word \Rightarrow int$
where

— treats the most-significant-bit as a sign bit

$sint-uint: sint\ w = sbintrunc\ (len-of\ TYPE\ ('a) - 1)\ (uint\ w)$

definition $unat :: 'a::len0\ word \Rightarrow nat$
where

$unat\ w = nat\ (uint\ w)$

definition $uints :: nat \Rightarrow int\ set$
where

— the sets of integers representing the words

$uints\ n = range\ (bintrunc\ n)$

definition $sints :: nat \Rightarrow int\ set$
where

$sints\ n = range\ (sbintrunc\ (n - 1))$

lemma *uints-num*:

$uints\ n = \{i. 0 \leq i \wedge i < 2 \wedge n\}$
 $\langle proof \rangle$

lemma *sints-num*:

$sints\ n = \{i. -(2 \wedge (n - 1)) \leq i \wedge i < 2 \wedge (n - 1)\}$
 $\langle proof \rangle$

definition *unats* :: *nat* \Rightarrow *nat set*

where

$unats\ n = \{i. i < 2 \wedge n\}$

definition *norm-sint* :: *nat* \Rightarrow *int* \Rightarrow *int*

where

$norm-sint\ n\ w = (w + 2 \wedge (n - 1)) \bmod 2 \wedge n - 2 \wedge (n - 1)$

definition *scast* :: '*a*::*len word* \Rightarrow '*b*::*len word*

where

— cast a word to a different length

$scast\ w = word-of-int\ (sint\ w)$

definition *ucast* :: '*a*::*len0 word* \Rightarrow '*b*::*len0 word*

where

$ucast\ w = word-of-int\ (uint\ w)$

instantiation *word* :: (*len0*) *size*

begin

definition

word-size: $size\ (w :: 'a\ word) = len-of\ TYPE('a)$

instance $\langle proof \rangle$

end

lemma *word-size-gt-0* [*iff*]:

$0 < size\ (w :: 'a :: len\ word)$
 $\langle proof \rangle$

lemmas *lens-gt-0* = *word-size-gt-0 len-gt-0*

lemma *lens-not-0* [*iff*]:

shows $size\ (w :: 'a :: len\ word) \neq 0$

and $len-of\ TYPE('a :: len) \neq 0$

$\langle proof \rangle$

definition *source-size* :: ('*a*::*len0 word* \Rightarrow '*b*) \Rightarrow *nat*

where

— whether a cast (or other) function is to a longer or shorter length
`[code del]: source-size c = (let arb = undefined; x = c arb in size arb)`

definition *target-size* :: ('a ⇒ 'b::len0 word) ⇒ nat
where
`[code del]: target-size c = size (c undefined)`

definition *is-up* :: ('a::len0 word ⇒ 'b::len0 word) ⇒ bool
where
`is-up c ⟷ source-size c ≤ target-size c`

definition *is-down* :: ('a :: len0 word ⇒ 'b :: len0 word) ⇒ bool
where
`is-down c ⟷ target-size c ≤ source-size c`

definition *of-bl* :: bool list ⇒ 'a::len0 word
where
`of-bl bl = word-of-int (bl-to-bin bl)`

definition *to-bl* :: 'a::len0 word ⇒ bool list
where
`to-bl w = bin-to-bl (len-of TYPE ('a)) (uint w)`

definition *word-reverse* :: 'a::len0 word ⇒ 'a word
where
`word-reverse w = of-bl (rev (to-bl w))`

definition *word-int-case* :: (int ⇒ 'b) ⇒ 'a::len0 word ⇒ 'b
where
`word-int-case f w = f (uint w)`

translations

`case x of XCONST of-int y => b == CONST word-int-case (%y. b) x`
`case x of (XCONST of-int :: 'a) y => b => CONST word-int-case (%y. b) x`

16.3 Correspondence relation for theorem transfer

definition *cr-word* :: int ⇒ 'a::len0 word ⇒ bool
where
`cr-word = (λx y. word-of-int x = y)`

lemma *Quotient-word*:

`Quotient (λx y. bintrunc (len-of TYPE('a)) x = bintrunc (len-of TYPE('a)) y)`
`word-of-int uint (cr-word :: - ⇒ 'a::len0 word ⇒ bool)`
`<proof>`

lemma *reflp-word*:

`reflp (λx y. bintrunc (len-of TYPE('a::len0)) x = bintrunc (len-of TYPE('a)) y)`

⟨proof⟩

setup-lifting *Quotient-word reflp-word*

TODO: The next lemma could be generated automatically.

lemma *uint-transfer* [*transfer-rule*]:

(*rel-fun pcr-word op =*) (*bintrunc (len-of TYPE('a))*)
 (*uint :: 'a::len0 word ⇒ int*)
 ⟨proof⟩

16.4 Basic code generation setup

definition *Word* :: *int ⇒ 'a::len0 word*

where

[*code-post*]: *Word = word-of-int*

lemma [*code abstype*]:

Word (uint w) = w
 ⟨proof⟩

declare *uint-word-of-int* [*code abstract*]

instantiation *word* :: (*len0*) *equal*

begin

definition *equal-word* :: '*a word ⇒ 'a word ⇒ bool*

where

equal-word k l ⇔ HOL.equal (uint k) (uint l)

instance ⟨proof⟩

end

notation *fcomp* (**infixl** $\circ > 60$)

notation *scomp* (**infixl** $\circ \rightarrow 60$)

instantiation *word* :: (*{len0, typerep}*) *random*

begin

definition

*random-word i = Random.range i $\circ \rightarrow$ ($\lambda k. \text{Pair} ($
 $\text{let } j = \text{word-of-int (int-of-integer (integer-of-natural } k)) :: 'a \text{ word}$
 $\text{in } (j, \lambda :: \text{unit. Code-Evaluation.term-of } j)))$*

instance ⟨proof⟩

end

no-notation *fcomp* (**infixl** $\circ > 60$)

no-notation *scomp* (**infixl** $\circ \rightarrow 60$)

16.5 Type-definition locale instantiations

lemmas *uint-0* = *uint-nonnegative*

lemmas *uint-lt* = *uint-bounded*

lemmas *uint-mod-same* = *uint-idem*

lemma *td-ext-uint*:

td-ext (*uint* :: 'a word \Rightarrow int) *word-of-int* (*uints* (*len-of TYPE*('a::len0)))
 ($\lambda w::\text{int}. w \bmod 2 \wedge \text{len-of TYPE}('a)$)
<proof>

interpretation *word-uint*:

td-ext *uint*::'a::len0 word \Rightarrow int
word-of-int
uints (*len-of TYPE*('a::len0))
 $\lambda w. w \bmod 2 \wedge \text{len-of TYPE}('a::\text{len0})$
<proof>

lemmas *td-uint* = *word-uint.td-thm*

lemmas *int-word-uint* = *word-uint.eq-norm*

lemma *td-ext-ubin*:

td-ext (*uint* :: 'a word \Rightarrow int) *word-of-int* (*uints* (*len-of TYPE*('a::len0)))
 (*bintrunc* (*len-of TYPE*('a)))
<proof>

interpretation *word-ubin*:

td-ext *uint*::'a::len0 word \Rightarrow int
word-of-int
uints (*len-of TYPE*('a::len0))
bintrunc (*len-of TYPE*('a::len0))
<proof>

16.6 Arithmetic operations

lift-definition *word-succ* :: 'a::len0 word \Rightarrow 'a word **is** $\lambda x. x + 1$

<proof>

lift-definition *word-pred* :: 'a::len0 word \Rightarrow 'a word **is** $\lambda x. x - 1$

<proof>

instantiation *word* :: (len0) {*neg-numeral*, *Divides.div*, *comm-monoid-mult*, *comm-ring*}
begin

lift-definition *zero-word* :: 'a word **is** 0 *<proof>*

lift-definition *one-word* :: 'a word **is** 1 *<proof>*

lift-definition *plus-word* :: 'a word \Rightarrow 'a word \Rightarrow 'a word is op +
 ⟨proof⟩

lift-definition *minus-word* :: 'a word \Rightarrow 'a word \Rightarrow 'a word is op -
 ⟨proof⟩

lift-definition *uminus-word* :: 'a word \Rightarrow 'a word is *uminus*
 ⟨proof⟩

lift-definition *times-word* :: 'a word \Rightarrow 'a word \Rightarrow 'a word is op *
 ⟨proof⟩

definition

word-div-def: $a \text{ div } b = \text{word-of-int } (\text{uint } a \text{ div uint } b)$

definition

word-mod-def: $a \text{ mod } b = \text{word-of-int } (\text{uint } a \text{ mod uint } b)$

instance

⟨proof⟩

end

Legacy theorems:

lemma *word-arith-wis* [code]: **shows**

word-add-def: $a + b = \text{word-of-int } (\text{uint } a + \text{uint } b)$ **and**
word-sub-wi: $a - b = \text{word-of-int } (\text{uint } a - \text{uint } b)$ **and**
word-mult-def: $a * b = \text{word-of-int } (\text{uint } a * \text{uint } b)$ **and**
word-minus-def: $- a = \text{word-of-int } (- \text{uint } a)$ **and**
word-succ-alt: $\text{word-succ } a = \text{word-of-int } (\text{uint } a + 1)$ **and**
word-pred-alt: $\text{word-pred } a = \text{word-of-int } (\text{uint } a - 1)$ **and**
word-0-wi: $0 = \text{word-of-int } 0$ **and**
word-1-wi: $1 = \text{word-of-int } 1$
 ⟨proof⟩

lemmas *ariths* =

bintr-ariths [THEN *word-ubin.norm-eq-iff* [THEN *iffD1*], folded *word-ubin.eq-norm*]

lemma *wi-homs*:

shows

wi-hom-add: $\text{word-of-int } a + \text{word-of-int } b = \text{word-of-int } (a + b)$ **and**
wi-hom-sub: $\text{word-of-int } a - \text{word-of-int } b = \text{word-of-int } (a - b)$ **and**
wi-hom-mult: $\text{word-of-int } a * \text{word-of-int } b = \text{word-of-int } (a * b)$ **and**
wi-hom-neg: $- \text{word-of-int } a = \text{word-of-int } (- a)$ **and**
wi-hom-succ: $\text{word-succ } (\text{word-of-int } a) = \text{word-of-int } (a + 1)$ **and**
wi-hom-pred: $\text{word-pred } (\text{word-of-int } a) = \text{word-of-int } (a - 1)$
 ⟨proof⟩

lemmas *wi-hom-syms* = *wi-homs* [symmetric]

lemmas *word-of-int-homs* = *wi-homs word-0-wi word-1-wi*

lemmas *word-of-int-hom-syms* = *word-of-int-homs [symmetric]*

instance *word* :: (*len*) *comm-ring-1*
 ⟨*proof*⟩

lemma *word-of-nat*: *of-nat n = word-of-int (int n)*
 ⟨*proof*⟩

lemma *word-of-int*: *of-int = word-of-int*
 ⟨*proof*⟩

definition *udvd* :: '*a*::*len word* => '*a*::*len word* => *bool* (**infixl** *udvd* 50)
where
a udvd b = (*EX n* >= 0. *uint b = n * uint a*)

16.7 Ordering

instantiation *word* :: (*len0*) *linorder*
begin

definition
word-le-def: *a ≤ b* ↔ *uint a ≤ uint b*

definition
word-less-def: *a < b* ↔ *uint a < uint b*

instance
 ⟨*proof*⟩

end

definition *word-sle* :: '*a* :: *len word* => '*a word* => *bool* ((-/ <=s -) [50, 51] 50)
where
a <=s b = (*sint a <= sint b*)

definition *word-sless* :: '*a* :: *len word* => '*a word* => *bool* ((-/ <s -) [50, 51] 50)
where
(x <s y) = (*x <=s y* & *x ~ = y*)

16.8 Bit-wise operations

instantiation *word* :: (*len0*) *bits*
begin

lift-definition *bitNOT-word* :: '*a word* => '*a word* **is** *bitNOT*
 ⟨*proof*⟩

lift-definition *bitAND-word* :: 'a word \Rightarrow 'a word \Rightarrow 'a word **is** *bitAND*
 ⟨proof⟩

lift-definition *bitOR-word* :: 'a word \Rightarrow 'a word \Rightarrow 'a word **is** *bitOR*
 ⟨proof⟩

lift-definition *bitXOR-word* :: 'a word \Rightarrow 'a word \Rightarrow 'a word **is** *bitXOR*
 ⟨proof⟩

definition

word-test-bit-def: $\text{test-bit } a = \text{bin-nth } (\text{uint } a)$

definition

word-set-bit-def: $\text{set-bit } a \ n \ x =$
word-of-int (*bin-sc* $n \ x$ (*uint* a))

definition

word-set-bits-def: (*BITS* $n. f \ n$) = *of-bl* (*bl-of-nth* (*len-of TYPE* ('a)) f)

definition

word-lsb-def: $\text{lsb } a \longleftrightarrow \text{bin-last } (\text{uint } a)$

definition *shifl1* :: 'a word \Rightarrow 'a word

where

shifl1 $w = \text{word-of-int } (\text{uint } w \ \text{BIT } \text{False})$

definition *shiftr1* :: 'a word \Rightarrow 'a word

where

— shift right as unsigned or as signed, ie logical or arithmetic

shiftr1 $w = \text{word-of-int } (\text{bin-rest } (\text{uint } w))$

definition

shifl-def: $w \ll n = (\text{shifl1 } \wedge \wedge n) \ w$

definition

shiftr-def: $w \gg n = (\text{shiftr1 } \wedge \wedge n) \ w$

instance ⟨proof⟩

end

lemma [*code*]: **shows**

word-not-def: $\text{NOT } (a::'a::\text{len0 } \text{word}) = \text{word-of-int } (\text{NOT } (\text{uint } a))$ **and**

word-and-def: $(a::'a \ \text{word}) \ \text{AND } b = \text{word-of-int } (\text{uint } a \ \text{AND } \text{uint } b)$ **and**

word-or-def: $(a::'a \ \text{word}) \ \text{OR } b = \text{word-of-int } (\text{uint } a \ \text{OR } \text{uint } b)$ **and**

word-xor-def: $(a::'a \ \text{word}) \ \text{XOR } b = \text{word-of-int } (\text{uint } a \ \text{XOR } \text{uint } b)$

⟨proof⟩

instantiation *word* :: (*len*) *bitss*

begin

definition

word-msb-def:

$msb\ a \longleftrightarrow bin\text{-}sign\ (sint\ a) = -1$

instance $\langle proof \rangle$

end

definition $setBit :: 'a :: len0\ word \Rightarrow nat \Rightarrow 'a\ word$

where

$setBit\ w\ n = set\text{-}bit\ w\ n\ True$

definition $clearBit :: 'a :: len0\ word \Rightarrow nat \Rightarrow 'a\ word$

where

$clearBit\ w\ n = set\text{-}bit\ w\ n\ False$

16.9 Shift operations

definition $sshiftr1 :: 'a :: len\ word \Rightarrow 'a\ word$

where

$sshiftr1\ w = word\text{-}of\text{-}int\ (bin\text{-}rest\ (sint\ w))$

definition $bshiftr1 :: bool \Rightarrow 'a :: len\ word \Rightarrow 'a\ word$

where

$bshiftr1\ b\ w = of\text{-}bl\ (b\ \# \text{butlast}\ (to\text{-}bl\ w))$

definition $sshiftr :: 'a :: len\ word \Rightarrow nat \Rightarrow 'a\ word$ (**infixl** $>>>$ 55)

where

$w\ >>>\ n = (sshiftr1\ \wedge\wedge\ n)\ w$

definition $mask :: nat \Rightarrow 'a::len\ word$

where

$mask\ n = (1\ <<\ n) - 1$

definition $revcast :: 'a :: len0\ word \Rightarrow 'b :: len0\ word$

where

$revcast\ w = of\text{-}bl\ (takefill\ False\ (len\text{-}of\ TYPE('b))\ (to\text{-}bl\ w))$

definition $slice1 :: nat \Rightarrow 'a :: len0\ word \Rightarrow 'b :: len0\ word$

where

$slice1\ n\ w = of\text{-}bl\ (takefill\ False\ n\ (to\text{-}bl\ w))$

definition $slice :: nat \Rightarrow 'a :: len0\ word \Rightarrow 'b :: len0\ word$

where

$slice\ n\ w = slice1\ (size\ w - n)\ w$

16.10 Rotation

definition *rotater1* :: 'a list => 'a list

where

rotater1 *ys* =
 (case *ys* of [] => [] | *x* # *xs* => last *ys* # butlast *ys*)

definition *rotater* :: nat => 'a list => 'a list

where

rotater *n* = *rotater1* ^^ *n*

definition *word-rotr* :: nat => 'a :: len0 word => 'a :: len0 word

where

word-rotr *n* *w* = of-bl (*rotater* *n* (to-bl *w*))

definition *word-rotl* :: nat => 'a :: len0 word => 'a :: len0 word

where

word-rotl *n* *w* = of-bl (*rotate* *n* (to-bl *w*))

definition *word-roti* :: int => 'a :: len0 word => 'a :: len0 word

where

word-roti *i* *w* = (if *i* >= 0 then *word-rotr* (nat *i*) *w*
 else *word-rotl* (nat (- *i*)) *w*)

16.11 Split and cat operations

definition *word-cat* :: 'a :: len0 word => 'b :: len0 word => 'c :: len0 word

where

word-cat *a* *b* = *word-of-int* (*bin-cat* (*uint* *a*) (*len-of TYPE* ('b)) (*uint* *b*))

definition *word-split* :: 'a :: len0 word => ('b :: len0 word) * ('c :: len0 word)

where

word-split *a* =
 (case *bin-split* (*len-of TYPE* ('c)) (*uint* *a*) of
 (*u*, *v*) => (*word-of-int* *u*, *word-of-int* *v*))

definition *word-rcat* :: 'a :: len0 word list => 'b :: len0 word

where

word-rcat *ws* =
word-of-int (*bin-rcat* (*len-of TYPE* ('a)) (*map uint* *ws*))

definition *word-rsplit* :: 'a :: len0 word => 'b :: len word list

where

word-rsplit *w* =
map word-of-int (*bin-rsplit* (*len-of TYPE* ('b)) (*len-of TYPE* ('a), *uint* *w*))

definition *max-word* :: 'a::len word — Largest representable machine integer.

where

max-word = *word-of-int* (2 ^ *len-of TYPE*('a) - 1)

lemmas *of-nth-def* = *word-set-bits-def*

16.12 Theorems about typedefs

lemma *sint-sbintrunc'*:

sint (*word-of-int* *bin* :: 'a word) =
 (*sbintrunc* (*len-of TYPE* ('a :: len) - 1) *bin*)
 ⟨*proof*⟩

lemma *uint-sint*:

uint *w* = *bintrunc* (*len-of TYPE*('a)) (*sint* (*w* :: 'a :: len word))
 ⟨*proof*⟩

lemma *bintr-uint*:

fixes *w* :: 'a::len0 word
shows *len-of TYPE*('a) ≤ *n* ⇒ *bintrunc* *n* (*uint* *w*) = *uint* *w*
 ⟨*proof*⟩

lemma *wi-bintr*:

len-of TYPE('a::len0) ≤ *n* ⇒
word-of-int (*bintrunc* *n* *w*) = (*word-of-int* *w* :: 'a word)
 ⟨*proof*⟩

lemma *td-ext-sbin*:

td-ext (*sint* :: 'a word ⇒ int) *word-of-int* (*sints* (*len-of TYPE*('a::len)))
 (*sbintrunc* (*len-of TYPE*('a) - 1))
 ⟨*proof*⟩

lemma *td-ext-sint*:

td-ext (*sint* :: 'a word ⇒ int) *word-of-int* (*sints* (*len-of TYPE*('a::len)))
 (λ*w*. (*w* + 2^{(*len-of TYPE*('a) - 1)}) mod 2^{*len-of TYPE*('a) - 2^{(*len-of TYPE*('a) - 1)}})
 ⟨*proof*⟩

interpretation *word-sint*:

td-ext *sint* :: 'a::len word => int
word-of-int
sints (*len-of TYPE*('a::len))
 %*w*. (*w* + 2^{(*len-of TYPE*('a::len) - 1)}) mod 2^{*len-of TYPE*('a::len) - 2^{(*len-of TYPE*('a::len) - 1)}})
 ⟨*proof*⟩

interpretation *word-sbin*:

td-ext *sint* :: 'a::len word => int
word-of-int
sints (*len-of TYPE*('a::len))
sbintrunc (*len-of TYPE*('a::len) - 1)
 ⟨*proof*⟩

lemmas *int-word-sint* = *td-ext-sint* [*THEN td-ext.eq-norm*]

lemmas *td-sint* = *word-sint.td*

lemma *to-bl-def'*:

(*to-bl* :: 'a :: len0 word => bool list) =
bin-to-bl (len-of TYPE('a)) o *uint*
 ⟨*proof*⟩

lemmas *word-reverse-no-def* [*simp*] = *word-reverse-def* [*of numeral w*] **for** *w*

lemma *uints-mod*: *uints n* = *range* ($\lambda w. w \bmod 2 \wedge n$)
 ⟨*proof*⟩

lemma *word-numeral-alt*:

numeral b = *word-of-int* (*numeral b*)
 ⟨*proof*⟩

declare *word-numeral-alt* [*symmetric, code-abbrev*]

lemma *word-neg-numeral-alt*:

$-$ *numeral b* = *word-of-int* ($-$ *numeral b*)
 ⟨*proof*⟩

declare *word-neg-numeral-alt* [*symmetric, code-abbrev*]

lemma *word-numeral-transfer* [*transfer-rule*]:

(*rel-fun op* = *pcr-word*) *numeral numeral*
 (*rel-fun op* = *pcr-word*) ($-$ *numeral*) ($-$ *numeral*)
 ⟨*proof*⟩

lemma *uint-bintrunc* [*simp*]:

uint (*numeral bin* :: 'a word) =
bintrunc (len-of TYPE ('a :: len0)) (*numeral bin*)
 ⟨*proof*⟩

lemma *uint-bintrunc-neg* [*simp*]: *uint* ($-$ *numeral bin* :: 'a word) =

bintrunc (len-of TYPE ('a :: len0)) ($-$ *numeral bin*)
 ⟨*proof*⟩

lemma *sint-sbintrunc* [*simp*]:

sint (*numeral bin* :: 'a word) =
sbintrunc (len-of TYPE ('a :: len) - 1) (*numeral bin*)
 ⟨*proof*⟩

lemma *sint-sbintrunc-neg* [*simp*]: *sint* ($-$ *numeral bin* :: 'a word) =

sbintrunc (len-of TYPE ('a :: len) - 1) ($-$ *numeral bin*)
 ⟨*proof*⟩

lemma *unat-bintrunc* [simp]:

$$\begin{aligned} \text{unat } (\text{numeral bin} :: 'a :: \text{len0 word}) &= \\ \text{nat } (\text{bintrunc } (\text{len-of TYPE('a)}) (\text{numeral bin})) & \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *unat-bintrunc-neg* [simp]:

$$\begin{aligned} \text{unat } (- \text{numeral bin} :: 'a :: \text{len0 word}) &= \\ \text{nat } (\text{bintrunc } (\text{len-of TYPE('a)}) (- \text{numeral bin})) & \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *size-0-eq*: $\text{size } (w :: 'a :: \text{len0 word}) = 0 \implies v = w$

$\langle \text{proof} \rangle$

lemma *uint-ge-0* [iff]: $0 \leq \text{uint } (x :: 'a :: \text{len0 word})$

$\langle \text{proof} \rangle$

lemma *uint-lt2p* [iff]: $\text{uint } (x :: 'a :: \text{len0 word}) < 2^{\text{len-of TYPE('a)}}$

$\langle \text{proof} \rangle$

lemma *sint-ge*: $-(2^{\text{len-of TYPE('a)} - 1}) \leq \text{sint } (x :: 'a :: \text{len word})$

$\langle \text{proof} \rangle$

lemma *sint-lt*: $\text{sint } (x :: 'a :: \text{len word}) < 2^{\text{len-of TYPE('a)} - 1}$

$\langle \text{proof} \rangle$

lemma *sign-uint-Pls* [simp]:

$$\begin{aligned} \text{bin-sign } (\text{uint } x) &= 0 \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *uint-m2p-neg*: $\text{uint } (x :: 'a :: \text{len0 word}) - 2^{\text{len-of TYPE('a)}} < 0$

$\langle \text{proof} \rangle$

lemma *uint-m2p-not-non-neg*:

$$\begin{aligned} \neg 0 \leq \text{uint } (x :: 'a :: \text{len0 word}) - 2^{\text{len-of TYPE('a)}} & \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *lt2p-lem*:

$$\begin{aligned} \text{len-of TYPE('a)} \leq n \implies \text{uint } (w :: 'a :: \text{len0 word}) &< 2^n \\ \langle \text{proof} \rangle \end{aligned}$$

lemma *uint-le-0-iff* [simp]: $\text{uint } x \leq 0 \iff \text{uint } x = 0$

$\langle \text{proof} \rangle$

lemma *uint-nat*: $\text{uint } w = \text{int } (\text{unat } w)$

$\langle \text{proof} \rangle$

lemma *uint-numeral*:

$$\text{uint } (\text{numeral } b :: 'a :: \text{len0 word}) = \text{numeral } b \bmod 2^{\text{len-of TYPE('a)}}$$

$\langle \text{proof} \rangle$

lemma *uint-neg-numeral*:

$\text{uint } (- \text{ numeral } b :: 'a :: \text{len0 word}) = - \text{ numeral } b \text{ mod } 2 ^ \text{len-of TYPE}('a)$
 $\langle \text{proof} \rangle$

lemma *unat-numeral*:

$\text{unat } (\text{ numeral } b :: 'a :: \text{len0 word}) = \text{ numeral } b \text{ mod } 2 ^ \text{len-of TYPE} ('a)$
 $\langle \text{proof} \rangle$

lemma *sint-numeral*: $\text{sint } (\text{ numeral } b :: 'a :: \text{len word}) = (\text{ numeral } b +$
 $2 ^ (\text{len-of TYPE}('a) - 1)) \text{ mod } 2 ^ \text{len-of TYPE}('a) -$
 $2 ^ (\text{len-of TYPE}('a) - 1)$
 $\langle \text{proof} \rangle$

lemma *word-of-int-0* [*simp, code-post*]:

$\text{word-of-int } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *word-of-int-1* [*simp, code-post*]:

$\text{word-of-int } 1 = 1$
 $\langle \text{proof} \rangle$

lemma *word-of-int-neg-1* [*simp*]: $\text{word-of-int } (- 1) = - 1$

$\langle \text{proof} \rangle$

lemma *word-of-int-numeral* [*simp*] :

$(\text{word-of-int } (\text{ numeral } bin) :: 'a :: \text{len0 word}) = (\text{ numeral } bin)$
 $\langle \text{proof} \rangle$

lemma *word-of-int-neg-numeral* [*simp*]:

$(\text{word-of-int } (- \text{ numeral } bin) :: 'a :: \text{len0 word}) = (- \text{ numeral } bin)$
 $\langle \text{proof} \rangle$

lemma *word-int-case-wi*:

$\text{word-int-case } f (\text{word-of-int } i :: 'b \text{ word}) =$
 $f (i \text{ mod } 2 ^ \text{len-of TYPE}('b :: \text{len0}))$
 $\langle \text{proof} \rangle$

lemma *word-int-split*:

$P (\text{word-int-case } f x) =$
 $(\text{ALL } i. x = (\text{word-of-int } i :: 'b :: \text{len0 word}) \ \&$
 $0 \leq i \ \& \ i < 2 ^ \text{len-of TYPE}('b) \ \longrightarrow P (f i))$
 $\langle \text{proof} \rangle$

lemma *word-int-split-asm*:

$P (\text{word-int-case } f x) =$
 $(\sim (\text{EX } n. x = (\text{word-of-int } n :: 'b :: \text{len0 word}) \ \&$
 $0 \leq n \ \& \ n < 2 ^ \text{len-of TYPE}('b :: \text{len0}) \ \& \sim P (f n)))$

<proof>

lemmas *uint-range'* = *word-uint.Rep* [*unfolded uints-num mem-Collect-eq*]

lemmas *sint-range'* = *word-sint.Rep* [*unfolded One-nat-def sints-num mem-Collect-eq*]

lemma *uint-range-size*: $0 \leq \text{uint } w \ \& \ \text{uint } w < 2^{\text{size } w}$

<proof>

lemma *sint-range-size*:

$-(2^{\text{size } w - \text{Suc } 0}) \leq \text{sint } w \ \& \ \text{sint } w < 2^{\text{size } w - \text{Suc } 0}$

<proof>

lemma *sint-above-size*: $2^{\text{size } (w::'a::\text{len } \text{word}) - 1} \leq x \implies \text{sint } w < x$

<proof>

lemma *sint-below-size*:

$x \leq -(2^{\text{size } (w::'a::\text{len } \text{word}) - 1}) \implies x \leq \text{sint } w$

<proof>

16.13 Testing bits

lemma *test-bit-eq-iff*: $(\text{test-bit } (u::'a::\text{len } 0 \ \text{word}) = \text{test-bit } v) = (u = v)$

<proof>

lemma *test-bit-size* [*rule-format*]: $(w::'a::\text{len } 0 \ \text{word}) !! n \dashrightarrow n < \text{size } w$

<proof>

lemma *word-eq-iff*:

fixes $x \ y :: 'a::\text{len } 0 \ \text{word}$

shows $x = y \longleftrightarrow (\forall n < \text{len-of } \text{TYPE}('a). \ x !! n = y !! n)$

<proof>

lemma *word-eqI* [*rule-format*]:

fixes $u :: 'a::\text{len } 0 \ \text{word}$

shows $(\text{ALL } n. \ n < \text{size } u \dashrightarrow u !! n = v !! n) \implies u = v$

<proof>

lemma *word-eqD*: $(u::'a::\text{len } 0 \ \text{word}) = v \implies u !! x = v !! x$

<proof>

lemma *test-bit-bin'*: $w !! n = (n < \text{size } w \ \& \ \text{bin-nth } (\text{uint } w) \ n)$

<proof>

lemmas *test-bit-bin* = *test-bit-bin'* [*unfolded word-size*]

lemma *bin-nth-uint-imp*:

$\text{bin-nth } (\text{uint } (w::'a::\text{len } 0 \ \text{word})) \ n \implies n < \text{len-of } \text{TYPE}('a)$

<proof>

lemma *bin-nth-sint*:

fixes $w :: 'a::len\ word$

shows $len-of\ TYPE('a) \leq n \implies$

$bin-nth\ (sint\ w)\ n = bin-nth\ (sint\ w)\ (len-of\ TYPE('a) - 1)$

<proof>

lemma *td-bl*:

type-definition $(to-bl :: 'a::len0\ word \Rightarrow\ bool\ list)$

of-bl

$\{bl.\ length\ bl = len-of\ TYPE('a)\}$

<proof>

interpretation *word-bl*:

type-definition $to-bl :: 'a::len0\ word \Rightarrow\ bool\ list$

of-bl

$\{bl.\ length\ bl = len-of\ TYPE('a::len0)\}$

<proof>

lemmas $word-bl-Rep' = word-bl.Rep\ [unfolded\ mem-Collect-eq,\ iff]$

lemma *word-size-bl*: $size\ w = size\ (to-bl\ w)$

<proof>

lemma *to-bl-use-of-bl*:

$(to-bl\ w = bl) = (w = of-bl\ bl \wedge length\ bl = length\ (to-bl\ w))$

<proof>

lemma *to-bl-word-rev*: $to-bl\ (word-reverse\ w) = rev\ (to-bl\ w)$

<proof>

lemma *word-rev-rev* [*simp*]: $word-reverse\ (word-reverse\ w) = w$

<proof>

lemma *word-rev-gal*: $word-reverse\ w = u \implies word-reverse\ u = w$

<proof>

lemma *word-rev-gal'*: $u = word-reverse\ w \implies w = word-reverse\ u$

<proof>

lemma *length-bl-gt-0* [*iff*]: $0 < length\ (to-bl\ (x::'a::len\ word))$

<proof>

lemma *bl-not-Nil* [*iff*]: $to-bl\ (x::'a::len\ word) \neq []$

<proof>

lemma *length-bl-neq-0* [*iff*]: $length\ (to-bl\ (x::'a::len\ word)) \neq 0$

<proof>

lemma *hd-bl-sign-sint*: $hd (to-bl w) = (bin-sign (sint w) = -1)$
 ⟨proof⟩

lemma *of-bl-drop'*:
 $lend = length bl - len-of TYPE ('a :: len0) \implies$
 $of-bl (drop lend bl) = (of-bl bl :: 'a word)$
 ⟨proof⟩

lemma *test-bit-of-bl*:
 $(of-bl bl :: 'a :: len0 word) !! n = (rev bl ! n \wedge n < len-of TYPE('a) \wedge n < length bl)$
 ⟨proof⟩

lemma *no-of-bl*:
 $(numeral bin :: 'a :: len0 word) = of-bl (bin-to-bl (len-of TYPE ('a)) (numeral bin))$
 ⟨proof⟩

lemma *uint-bl*: $to-bl w = bin-to-bl (size w) (uint w)$
 ⟨proof⟩

lemma *to-bl-bin*: $bl-to-bin (to-bl w) = uint w$
 ⟨proof⟩

lemma *to-bl-of-bin*:
 $to-bl (word-of-int bin :: 'a :: len0 word) = bin-to-bl (len-of TYPE('a)) bin$
 ⟨proof⟩

lemma *to-bl-numeral [simp]*:
 $to-bl (numeral bin :: 'a :: len0 word) =$
 $bin-to-bl (len-of TYPE('a)) (numeral bin)$
 ⟨proof⟩

lemma *to-bl-neg-numeral [simp]*:
 $to-bl (- numeral bin :: 'a :: len0 word) =$
 $bin-to-bl (len-of TYPE('a)) (- numeral bin)$
 ⟨proof⟩

lemma *to-bl-to-bin [simp]*: $bl-to-bin (to-bl w) = uint w$
 ⟨proof⟩

lemma *uint-bl-bin*:
fixes $x :: 'a :: len0 word$
shows $bl-to-bin (bin-to-bl (len-of TYPE('a)) (uint x)) = uint x$
 ⟨proof⟩

lemma *uints-unats*: $uints n = int ' unats n$
 ⟨proof⟩

lemma *unats-wints*: $unats\ n = nat\ 'uints\ n$
 ⟨*proof*⟩

lemmas *bintr-num = word-ubin.norm-eq-iff*
 [of numeral a numeral b, symmetric, folded word-numeral-alt] **for** a b
lemmas *sbintr-num = word-sbin.norm-eq-iff*
 [of numeral a numeral b, symmetric, folded word-numeral-alt] **for** a b

lemma *num-of-bintr'*:
 $bintrunc\ (len-of\ TYPE('a :: len0))\ (numeral\ a) = (numeral\ b) \implies$
 $numeral\ a = (numeral\ b :: 'a\ word)$
 ⟨*proof*⟩

lemma *num-of-sbintr'*:
 $sbintrunc\ (len-of\ TYPE('a :: len) - 1)\ (numeral\ a) = (numeral\ b) \implies$
 $numeral\ a = (numeral\ b :: 'a\ word)$
 ⟨*proof*⟩

lemma *num-abs-bintr*:
 $(numeral\ x :: 'a\ word) =$
 $word-of-int\ (bintrunc\ (len-of\ TYPE('a::len0))\ (numeral\ x))$
 ⟨*proof*⟩

lemma *num-abs-sbintr*:
 $(numeral\ x :: 'a\ word) =$
 $word-of-int\ (sbintrunc\ (len-of\ TYPE('a::len) - 1)\ (numeral\ x))$
 ⟨*proof*⟩

lemma *ucast-id*: $ucast\ w = w$
 ⟨*proof*⟩

lemma *scast-id*: $scast\ w = w$
 ⟨*proof*⟩

lemma *ucast-bl*: $ucast\ w = of-bl\ (to-bl\ w)$
 ⟨*proof*⟩

lemma *nth-ucast*:
 $(ucast\ w :: 'a :: len0\ word) !! n = (w !! n \& n < len-of\ TYPE('a))$
 ⟨*proof*⟩

lemma *ucast-bintr [simp]*:
 $ucast\ (numeral\ w :: 'a :: len0\ word) =$
 $word-of-int\ (bintrunc\ (len-of\ TYPE('a))\ (numeral\ w))$
 ⟨*proof*⟩

lemma *scast-sbintr* [*simp*]:
 $scast (numeral w :: 'a::len0 word) =$
 $word-of-int (sbintrunc (len-of TYPE('a) - Suc 0) (numeral w))$
 ⟨*proof*⟩

lemma *source-size*: $source-size (c::'a::len0 word \Rightarrow -) = len-of TYPE('a)$
 ⟨*proof*⟩

lemma *target-size*: $target-size (c::- \Rightarrow 'b::len0 word) = len-of TYPE('b)$
 ⟨*proof*⟩

lemma *is-down*:
fixes $c :: 'a::len0 word \Rightarrow 'b::len0 word$
shows $is-down c \longleftrightarrow len-of TYPE('b) \leq len-of TYPE('a)$
 ⟨*proof*⟩

lemma *is-up*:
fixes $c :: 'a::len0 word \Rightarrow 'b::len0 word$
shows $is-up c \longleftrightarrow len-of TYPE('a) \leq len-of TYPE('b)$
 ⟨*proof*⟩

lemmas *is-up-down = trans* [*OF is-up is-down [symmetric]*]

lemma *down-cast-same* [*OF refl*]: $uc = ucast \Longrightarrow is-down uc \Longrightarrow uc = scast$
 ⟨*proof*⟩

lemma *word-rev-tf*:
 $to-bl (of-bl bl::'a::len0 word) =$
 $rev (takefill False (len-of TYPE('a)) (rev bl))$
 ⟨*proof*⟩

lemma *word-rep-drop*:
 $to-bl (of-bl bl::'a::len0 word) =$
 $replicate (len-of TYPE('a) - length bl) False @$
 $drop (length bl - len-of TYPE('a)) bl$
 ⟨*proof*⟩

lemma *to-bl-ucast*:
 $to-bl (ucast (w::'b::len0 word) ::'a::len0 word) =$
 $replicate (len-of TYPE('a) - len-of TYPE('b)) False @$
 $drop (len-of TYPE('b) - len-of TYPE('a)) (to-bl w)$
 ⟨*proof*⟩

lemma *ucast-up-app* [*OF refl*]:
 $uc = ucast \Longrightarrow source-size uc + n = target-size uc \Longrightarrow$
 $to-bl (uc w) = replicate n False @ (to-bl w)$
 ⟨*proof*⟩

lemma *ucast-down-drop* [*OF refl*]:

$$uc = ucast \implies source\text{-}size\ uc = target\text{-}size\ uc + n \implies$$

$$to\text{-}bl\ (uc\ w) = drop\ n\ (to\text{-}bl\ w)$$

<proof>

lemma *scast-down-drop* [*OF refl*]:

$$sc = scast \implies source\text{-}size\ sc = target\text{-}size\ sc + n \implies$$

$$to\text{-}bl\ (sc\ w) = drop\ n\ (to\text{-}bl\ w)$$

<proof>

lemma *sint-up-scast* [*OF refl*]:

$$sc = scast \implies is\text{-}up\ sc \implies sint\ (sc\ w) = sint\ w$$

<proof>

lemma *uint-up-ucast* [*OF refl*]:

$$uc = ucast \implies is\text{-}up\ uc \implies uint\ (uc\ w) = uint\ w$$

<proof>

lemma *ucast-up-ucast* [*OF refl*]:

$$uc = ucast \implies is\text{-}up\ uc \implies ucast\ (uc\ w) = ucast\ w$$

<proof>

lemma *scast-up-scast* [*OF refl*]:

$$sc = scast \implies is\text{-}up\ sc \implies scast\ (sc\ w) = scast\ w$$

<proof>

lemma *ucast-of-bl-up* [*OF refl*]:

$$w = of\text{-}bl\ bl \implies size\ bl \leq size\ w \implies ucast\ w = of\text{-}bl\ bl$$

<proof>

lemmas *ucast-up-ucast-id = trans* [*OF ucast-up-ucast ucast-id*]

lemmas *scast-up-scast-id = trans* [*OF scast-up-scast scast-id*]

lemmas *isdou = is-up-down* [**where** *c = ucast*, *THEN iffD2*]

lemmas *isdus = is-up-down* [**where** *c = scast*, *THEN iffD2*]

lemmas *ucast-down-ucast-id = isdou* [*THEN ucast-up-ucast-id*]

lemmas *scast-down-scast-id = isdus* [*THEN ucast-up-ucast-id*]

lemma *up-ucast-surj*:

$$is\text{-}up\ (ucast\ ::\ 'b::len0\ word \implies 'a::len0\ word) \implies$$

$$surj\ (ucast\ ::\ 'a\ word \implies 'b\ word)$$

<proof>

lemma *up-scast-surj*:

$$is\text{-}up\ (scast\ ::\ 'b::len\ word \implies 'a::len\ word) \implies$$

$$surj\ (scast\ ::\ 'a\ word \implies 'b\ word)$$

<proof>

lemma *down-scast-inj*:

is-down (*scast* :: 'b::len word => 'a::len word) \implies
inj-on (*ucast* :: 'a word => 'b word) *A*
 ⟨*proof*⟩

lemma *down-ucast-inj*:

is-down (*ucast* :: 'b::len0 word => 'a::len0 word) \implies
inj-on (*ucast* :: 'a word => 'b word) *A*
 ⟨*proof*⟩

lemma *of-bl-append-same*: *of-bl* (*X* @ *to-bl w*) = *w*

⟨*proof*⟩

lemma *ucast-down-wi* [*OF refl*]:

uc = *ucast* \implies *is-down uc* \implies *uc* (*word-of-int x*) = *word-of-int x*
 ⟨*proof*⟩

lemma *ucast-down-no* [*OF refl*]:

uc = *ucast* \implies *is-down uc* \implies *uc* (*numeral bin*) = *numeral bin*
 ⟨*proof*⟩

lemma *ucast-down-bl* [*OF refl*]:

uc = *ucast* \implies *is-down uc* \implies *uc* (*of-bl bl*) = *of-bl bl*
 ⟨*proof*⟩

lemmas *slice-def'* = *slice-def* [*unfolded word-size*]

lemmas *test-bit-def'* = *word-test-bit-def* [*THEN fun-cong*]

lemmas *word-log-defs* = *word-and-def* *word-or-def* *word-xor-def* *word-not-def*

16.14 Word Arithmetic

lemma *word-less-alt*: (*a* < *b*) = (*uint a* < *uint b*)

⟨*proof*⟩

lemma *signed-linorder*: *class.linorder* *word-sle* *word-sless*

⟨*proof*⟩

interpretation *signed*: *linorder* *word-sle* *word-sless*

⟨*proof*⟩

lemma *udvdI*:

$0 \leq n \implies$ *uint b* = *n* * *uint a* \implies *a* *udvd* *b*
 ⟨*proof*⟩

lemmas *word-div-no* [*simp*] = *word-div-def* [*of numeral a numeral b*] **for** *a b*

lemmas *word-mod-no* [*simp*] = *word-mod-def* [*of numeral a numeral b*] **for** *a b*

lemmas *word-less-no* [simp] = *word-less-def* [of numeral a numeral b] **for** a b

lemmas *word-le-no* [simp] = *word-le-def* [of numeral a numeral b] **for** a b

lemmas *word-sless-no* [simp] = *word-sless-def* [of numeral a numeral b] **for** a b

lemmas *word-sle-no* [simp] = *word-sle-def* [of numeral a numeral b] **for** a b

lemma *word-m1-wi*: $- 1 = \text{word-of-int } (- 1)$
 ⟨proof⟩

lemma *word-0-bl* [simp]: *of-bl* [] = 0
 ⟨proof⟩

lemma *word-1-bl*: *of-bl* [True] = 1
 ⟨proof⟩

lemma *uint-eq-0* [simp]: *uint* 0 = 0
 ⟨proof⟩

lemma *of-bl-0* [simp]: *of-bl* (replicate n False) = 0
 ⟨proof⟩

lemma *to-bl-0* [simp]:
to-bl (0 :: 'a :: len0 word) = replicate (len-of TYPE('a)) False
 ⟨proof⟩

lemma *uint-0-iff*:
uint x = 0 \longleftrightarrow x = 0
 ⟨proof⟩

lemma *unat-0-iff*:
unat x = 0 \longleftrightarrow x = 0
 ⟨proof⟩

lemma *unat-0* [simp]:
unat 0 = 0
 ⟨proof⟩

lemma *size-0-same'*:
size w = 0 \implies w = (v :: 'a :: len0 word)
 ⟨proof⟩

lemmas *size-0-same* = *size-0-same'* [unfolded word-size]

lemmas *unat-eq-0* = *unat-0-iff*

lemmas *unat-eq-zero* = *unat-0-iff*

lemma *unat-gt-0*: (0 < *unat* x) = (x \sim = 0)

<proof>

lemma *ucast-0* [*simp*]: *ucast 0 = 0*
<proof>

lemma *sint-0* [*simp*]: *sint 0 = 0*
<proof>

lemma *scast-0* [*simp*]: *scast 0 = 0*
<proof>

lemma *sint-n1* [*simp*]: *sint (- 1) = - 1*
<proof>

lemma *scast-n1* [*simp*]: *scast (- 1) = - 1*
<proof>

lemma *uint-1* [*simp*]: *uint (1::'a::len word) = 1*
<proof>

lemma *unat-1* [*simp*]: *unat (1::'a::len word) = 1*
<proof>

lemma *ucast-1* [*simp*]: *ucast (1::'a::len word) = 1*
<proof>

16.15 Transferring goals from words to ints

lemma *word-ths*:

shows

word-succ-p1: *word-succ a = a + 1* **and**

word-pred-m1: *word-pred a = a - 1* **and**

word-pred-succ: *word-pred (word-succ a) = a* **and**

word-succ-pred: *word-succ (word-pred a) = a* **and**

word-mult-succ: *word-succ a * b = b + a * b*

<proof>

lemma *uint-cong*: *x = y \implies uint x = uint y*
<proof>

lemma *uint-word-ariths*:

fixes *a b :: 'a::len0 word*

shows *uint (a + b) = (uint a + uint b) mod 2 ^ len-of TYPE('a::len0)*

and *uint (a - b) = (uint a - uint b) mod 2 ^ len-of TYPE('a)*

and *uint (a * b) = uint a * uint b mod 2 ^ len-of TYPE('a)*

and *uint (- a) = - uint a mod 2 ^ len-of TYPE('a)*

and *uint (word-succ a) = (uint a + 1) mod 2 ^ len-of TYPE('a)*

and *uint (word-pred a) = (uint a - 1) mod 2 ^ len-of TYPE('a)*

and *uint (0 :: 'a word) = 0 mod 2 ^ len-of TYPE('a)*

and $\text{uint } (1 :: 'a \text{ word}) = 1 \text{ mod } 2 \wedge \text{len-of TYPE}('a)$
 ⟨proof⟩

lemma *uint-word-arith-bintrs*:

fixes $a \ b :: 'a::\text{len0 word}$
shows $\text{uint } (a + b) = \text{bintrunc } (\text{len-of TYPE}('a)) (\text{uint } a + \text{uint } b)$
and $\text{uint } (a - b) = \text{bintrunc } (\text{len-of TYPE}('a)) (\text{uint } a - \text{uint } b)$
and $\text{uint } (a * b) = \text{bintrunc } (\text{len-of TYPE}('a)) (\text{uint } a * \text{uint } b)$
and $\text{uint } (- a) = \text{bintrunc } (\text{len-of TYPE}('a)) (- \text{uint } a)$
and $\text{uint } (\text{word-succ } a) = \text{bintrunc } (\text{len-of TYPE}('a)) (\text{uint } a + 1)$
and $\text{uint } (\text{word-pred } a) = \text{bintrunc } (\text{len-of TYPE}('a)) (\text{uint } a - 1)$
and $\text{uint } (0 :: 'a \text{ word}) = \text{bintrunc } (\text{len-of TYPE}('a)) 0$
and $\text{uint } (1 :: 'a \text{ word}) = \text{bintrunc } (\text{len-of TYPE}('a)) 1$
 ⟨proof⟩

lemma *sint-word-ariths*:

fixes $a \ b :: 'a::\text{len word}$
shows $\text{sint } (a + b) = \text{sbintrunc } (\text{len-of TYPE}('a) - 1) (\text{sint } a + \text{sint } b)$
and $\text{sint } (a - b) = \text{sbintrunc } (\text{len-of TYPE}('a) - 1) (\text{sint } a - \text{sint } b)$
and $\text{sint } (a * b) = \text{sbintrunc } (\text{len-of TYPE}('a) - 1) (\text{sint } a * \text{sint } b)$
and $\text{sint } (- a) = \text{sbintrunc } (\text{len-of TYPE}('a) - 1) (- \text{sint } a)$
and $\text{sint } (\text{word-succ } a) = \text{sbintrunc } (\text{len-of TYPE}('a) - 1) (\text{sint } a + 1)$
and $\text{sint } (\text{word-pred } a) = \text{sbintrunc } (\text{len-of TYPE}('a) - 1) (\text{sint } a - 1)$
and $\text{sint } (0 :: 'a \text{ word}) = \text{sbintrunc } (\text{len-of TYPE}('a) - 1) 0$
and $\text{sint } (1 :: 'a \text{ word}) = \text{sbintrunc } (\text{len-of TYPE}('a) - 1) 1$
 ⟨proof⟩

lemmas $\text{uint-div-alt} = \text{word-div-def } [\text{THEN trans } [\text{OF uint-cong int-word-uint}]]$

lemmas $\text{uint-mod-alt} = \text{word-mod-def } [\text{THEN trans } [\text{OF uint-cong int-word-uint}]]$

lemma *word-pred-0-n1*: $\text{word-pred } 0 = \text{word-of-int } (- 1)$
 ⟨proof⟩

lemma *succ-pred-no* [simp]:

$\text{word-succ } (\text{numeral } w) = \text{numeral } w + 1$
 $\text{word-pred } (\text{numeral } w) = \text{numeral } w - 1$
 $\text{word-succ } (- \text{numeral } w) = - \text{numeral } w + 1$
 $\text{word-pred } (- \text{numeral } w) = - \text{numeral } w - 1$
 ⟨proof⟩

lemma *word-sp-01* [simp] :

$\text{word-succ } (- 1) = 0 \ \& \ \text{word-succ } 0 = 1 \ \& \ \text{word-pred } 0 = - 1 \ \& \ \text{word-pred } 1 = 0$
 ⟨proof⟩

lemma *word-of-int-Ex*:

$\exists y. x = \text{word-of-int } y$
 ⟨proof⟩

16.16 Order on fixed-length words

lemma *word-zero-le* [*simp*] :
 $0 \leq (y :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemma *word-m1-ge* [*simp*] : *word-pred* 0 $\geq y$
 ⟨*proof*⟩

lemma *word-n1-ge* [*simp*]: $y \leq (-1 :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemmas *word-not-simps* [*simp*] =
word-zero-le [THEN *leD*] *word-m1-ge* [THEN *leD*] *word-n1-ge* [THEN *leD*]

lemma *word-gt-0*: $0 < y \longleftrightarrow 0 \neq (y :: 'a :: \text{len0 word})$
 ⟨*proof*⟩

lemmas *word-gt-0-no* [*simp*] = *word-gt-0* [*of numeral* *y*] **for** *y*

lemma *word-sless-alt*: $(a <_s b) = (\text{sint } a < \text{sint } b)$
 ⟨*proof*⟩

lemma *word-le-nat-alt*: $(a \leq b) = (\text{unat } a \leq \text{unat } b)$
 ⟨*proof*⟩

lemma *word-less-nat-alt*: $(a < b) = (\text{unat } a < \text{unat } b)$
 ⟨*proof*⟩

lemma *wi-less*:
 $(\text{word-of-int } n < (\text{word-of-int } m :: 'a :: \text{len0 word})) =$
 $(n \bmod 2 \wedge \text{len-of TYPE('a)} < m \bmod 2 \wedge \text{len-of TYPE('a)})$
 ⟨*proof*⟩

lemma *wi-le*:
 $(\text{word-of-int } n \leq (\text{word-of-int } m :: 'a :: \text{len0 word})) =$
 $(n \bmod 2 \wedge \text{len-of TYPE('a)} \leq m \bmod 2 \wedge \text{len-of TYPE('a)})$
 ⟨*proof*⟩

lemma *udvd-nat-alt*: $a \text{ udvd } b = (\exists n \geq 0. \text{unat } b = n * \text{unat } a)$
 ⟨*proof*⟩

lemma *udvd-iff-dvd*: $x \text{ udvd } y \longleftrightarrow \text{unat } x \text{ dvd } \text{unat } y$
 ⟨*proof*⟩

lemmas *unat-mono* = *word-less-nat-alt* [THEN *iffD1*]

lemma *unat-minus-one*:
assumes $w \neq 0$
shows $\text{unat } (w - 1) = \text{unat } w - 1$

\langle proof \rangle

lemma *measure-unat*: $p \approx 0 \implies \text{unat } (p - 1) < \text{unat } p$
 \langle proof \rangle

lemmas *uint-add-ge0* [simp] = *add-nonneg-nonneg* [OF *uint-ge-0 uint-ge-0*]

lemmas *uint-mult-ge0* [simp] = *mult-nonneg-nonneg* [OF *uint-ge-0 uint-ge-0*]

lemma *uint-sub-lt2p* [simp]:
 $\text{uint } (x :: 'a :: \text{len0 word}) - \text{uint } (y :: 'b :: \text{len0 word}) <$
 $2 \wedge \text{len-of TYPE}('a)$
 \langle proof \rangle

16.17 Conditions for the addition (etc) of two words to overflow

lemma *uint-add-lem*:
 $(\text{uint } x + \text{uint } y < 2 \wedge \text{len-of TYPE}('a)) =$
 $(\text{uint } (x + y :: 'a :: \text{len0 word}) = \text{uint } x + \text{uint } y)$
 \langle proof \rangle

lemma *uint-mult-lem*:
 $(\text{uint } x * \text{uint } y < 2 \wedge \text{len-of TYPE}('a)) =$
 $(\text{uint } (x * y :: 'a :: \text{len0 word}) = \text{uint } x * \text{uint } y)$
 \langle proof \rangle

lemma *uint-sub-lem*:
 $(\text{uint } x \geq \text{uint } y) = (\text{uint } (x - y) = \text{uint } x - \text{uint } y)$
 \langle proof \rangle

lemma *uint-add-le*: $\text{uint } (x + y) \leq \text{uint } x + \text{uint } y$
 \langle proof \rangle

lemma *uint-sub-ge*: $\text{uint } (x - y) \geq \text{uint } x - \text{uint } y$
 \langle proof \rangle

lemma *mod-add-if-z*:
 $(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$
 $(x + y) \bmod z = (\text{if } x + y < z \text{ then } x + y \text{ else } x + y - z)$
 \langle proof \rangle

lemma *uint-plus-if'*:
 $\text{uint } ((a :: 'a \text{ word}) + b) =$
 $(\text{if } \text{uint } a + \text{uint } b < 2 \wedge \text{len-of TYPE}('a :: \text{len0}) \text{ then } \text{uint } a + \text{uint } b$
 $\text{else } \text{uint } a + \text{uint } b - 2 \wedge \text{len-of TYPE}('a))$
 \langle proof \rangle

lemma *mod-sub-if-z*:
 $(x :: \text{int}) < z \implies y < z \implies 0 \leq y \implies 0 \leq x \implies 0 \leq z \implies$

$(x - y) \bmod z = (\text{if } y \leq x \text{ then } x - y \text{ else } x - y + z)$
 ⟨proof⟩

lemma *uint-sub-if'*:

$\text{uint } ((a :: 'a \text{ word}) - b) =$
 (if $\text{uint } b \leq \text{uint } a$ then $\text{uint } a - \text{uint } b$
 else $\text{uint } a - \text{uint } b + 2^{\wedge} \text{len-of TYPE}('a :: \text{len}0)$)
 ⟨proof⟩

16.18 Definition of *uint-arith*

lemma *word-of-int-inverse*:

$\text{word-of-int } r = a \implies 0 \leq r \implies r < 2^{\wedge} \text{len-of TYPE}('a) \implies$
 $\text{uint } (a :: 'a :: \text{len}0 \text{ word}) = r$
 ⟨proof⟩

lemma *uint-split*:

fixes $x :: 'a :: \text{len}0 \text{ word}$
shows $P (\text{uint } x) =$
 (ALL $i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2^{\wedge} \text{len-of TYPE}('a) \implies P i)$
 ⟨proof⟩

lemma *uint-split-asm*:

fixes $x :: 'a :: \text{len}0 \text{ word}$
shows $P (\text{uint } x) =$
 ($\sim (EX i. \text{word-of-int } i = x \ \& \ 0 \leq i \ \& \ i < 2^{\wedge} \text{len-of TYPE}('a) \ \& \ \sim P i)$)
 ⟨proof⟩

lemmas *uint-splits = uint-split uint-split-asm*

lemmas *uint-arith-simps =*

word-le-def word-less-alt
word-uint.Rep-inject [symmetric]
uint-sub-if' uint-plus-if'

lemma *power-False-cong*: $\text{False} \implies a^{\wedge} b = c^{\wedge} d$
 ⟨proof⟩

⟨ML⟩

16.19 More on overflows and monotonicity

lemma *no-plus-overflow-uint-size*:

$((x :: 'a :: \text{len}0 \text{ word}) \leq x + y) = (\text{uint } x + \text{uint } y < 2^{\wedge} \text{size } x)$
 ⟨proof⟩

lemmas *no-olen-add = no-plus-overflow-uint-size [unfolded word-size]*

lemma *no-ulen-sub*: $((x :: 'a :: \text{len0 word}) \geq x - y) = (\text{uint } y \leq \text{uint } x)$
 ⟨proof⟩

lemma *no-olen-add'*:
fixes $x :: 'a :: \text{len0 word}$
shows $(x \leq y + x) = (\text{uint } y + \text{uint } x < 2 \wedge \text{len-of TYPE('a)})$
 ⟨proof⟩

lemmas *olen-add-eqv* = *trans* [*OF no-olen-add no-olen-add'* [*symmetric*]]

lemmas *uint-plus-simple-iff* = *trans* [*OF no-olen-add uint-add-lem*]
lemmas *uint-plus-simple* = *uint-plus-simple-iff* [*THEN iffD1*]
lemmas *uint-minus-simple-iff* = *trans* [*OF no-ulen-sub uint-sub-lem*]
lemmas *uint-minus-simple-alt* = *uint-sub-lem* [*folded word-le-def*]
lemmas *word-sub-le-iff* = *no-ulen-sub* [*folded word-le-def*]
lemmas *word-sub-le* = *word-sub-le-iff* [*THEN iffD2*]

lemma *word-less-sub1*:
 $(x :: 'a :: \text{len word}) \sim 0 \implies (1 < x) = (0 < x - 1)$
 ⟨proof⟩

lemma *word-le-sub1*:
 $(x :: 'a :: \text{len word}) \sim 0 \implies (1 \leq x) = (0 \leq x - 1)$
 ⟨proof⟩

lemma *sub-wrap-lt*:
 $((x :: 'a :: \text{len0 word}) < x - z) = (x < z)$
 ⟨proof⟩

lemma *sub-wrap*:
 $((x :: 'a :: \text{len0 word}) \leq x - z) = (z = 0 \mid x < z)$
 ⟨proof⟩

lemma *plus-minus-not-NULL-ab*:
 $(x :: 'a :: \text{len0 word}) \leq ab - c \implies c \leq ab \implies c \sim 0 \implies x + c \sim 0$
 ⟨proof⟩

lemma *plus-minus-no-overflow-ab*:
 $(x :: 'a :: \text{len0 word}) \leq ab - c \implies c \leq ab \implies x \leq x + c$
 ⟨proof⟩

lemma *le-minus'*:
 $(a :: 'a :: \text{len0 word}) + c \leq b \implies a \leq a + c \implies c \leq b - a$
 ⟨proof⟩

lemma *le-plus'*:
 $(a :: 'a :: \text{len0 word}) \leq b \implies c \leq b - a \implies a + c \leq b$
 ⟨proof⟩

lemmas *le-plus = le-plus'* [rotated]

lemmas *le-minus = leD* [THEN *thin-rl*, THEN *le-minus'*]

lemma *word-plus-mono-right:*

$(y :: 'a :: \text{len0 word}) <= z \implies x <= x + z \implies x + y <= x + z$
 ⟨proof⟩

lemma *word-less-minus-cancel:*

$y - x < z - x \implies x <= z \implies (y :: 'a :: \text{len0 word}) < z$
 ⟨proof⟩

lemma *word-less-minus-mono-left:*

$(y :: 'a :: \text{len0 word}) < z \implies x <= y \implies y - x < z - x$
 ⟨proof⟩

lemma *word-less-minus-mono:*

$a < c \implies d < b \implies a - b < a \implies c - d < c$
 $\implies a - b < c - (d :: 'a :: \text{len word})$
 ⟨proof⟩

lemma *word-le-minus-cancel:*

$y - x <= z - x \implies x <= z \implies (y :: 'a :: \text{len0 word}) <= z$
 ⟨proof⟩

lemma *word-le-minus-mono-left:*

$(y :: 'a :: \text{len0 word}) <= z \implies x <= y \implies y - x <= z - x$
 ⟨proof⟩

lemma *word-le-minus-mono:*

$a <= c \implies d <= b \implies a - b <= a \implies c - d <= c$
 $\implies a - b <= c - (d :: 'a :: \text{len word})$
 ⟨proof⟩

lemma *plus-le-left-cancel-wrap:*

$(x :: 'a :: \text{len0 word}) + y' < x \implies x + y < x \implies (x + y' < x + y) = (y' < y)$
 ⟨proof⟩

lemma *plus-le-left-cancel-nowrap:*

$(x :: 'a :: \text{len0 word}) <= x + y' \implies x <= x + y \implies$
 $(x + y' < x + y) = (y' < y)$
 ⟨proof⟩

lemma *word-plus-mono-right2:*

$(a :: 'a :: \text{len0 word}) <= a + b \implies c <= b \implies a <= a + c$
 ⟨proof⟩

lemma *word-less-add-right:*

$(x :: 'a :: \text{len0 word}) < y - z \implies z <= y \implies x + z < y$

<proof>

lemma *word-less-sub-right*:

$(x :: 'a :: \text{len0 word}) < y + z \implies y \leq x \implies x - y < z$
<proof>

lemma *word-le-plus-either*:

$(x :: 'a :: \text{len0 word}) \leq y \mid x \leq z \implies y \leq y + z \implies x \leq y + z$
<proof>

lemma *word-less-nowrapI*:

$(x :: 'a :: \text{len0 word}) < z - k \implies k \leq z \implies 0 < k \implies x < x + k$
<proof>

lemma *inc-le*: $(i :: 'a :: \text{len word}) < m \implies i + 1 \leq m$

<proof>

lemma *inc-i*:

$(1 :: 'a :: \text{len word}) \leq i \implies i < m \implies 1 \leq (i + 1) \ \& \ i + 1 \leq m$
<proof>

lemma *udvd-incr-lem*:

$up < uq \implies up = ua + n * \text{uint } K \implies$
 $uq = ua + n' * \text{uint } K \implies up + \text{uint } K \leq uq$
<proof>

lemma *udvd-incr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p + K \leq q$
<proof>

lemma *udvd-decr'*:

$p < q \implies \text{uint } p = ua + n * \text{uint } K \implies$
 $\text{uint } q = ua + n' * \text{uint } K \implies p \leq q - K$
<proof>

lemmas *udvd-incr-lem0* = *udvd-incr-lem* [**where** *ua=0, unfolded add-0-left*]

lemmas *udvd-incr0* = *udvd-incr'* [**where** *ua=0, unfolded add-0-left*]

lemmas *udvd-decr0* = *udvd-decr'* [**where** *ua=0, unfolded add-0-left*]

lemma *udvd-minus-le'*:

$xy < k \implies z \text{ udvd } xy \implies z \text{ udvd } k \implies xy \leq k - z$
<proof>

lemma *udvd-incr2-K*:

$p < a + s \implies a \leq a + s \implies K \text{ udvd } s \implies K \text{ udvd } p - a \implies a \leq p \implies$
 $0 < K \implies p \leq p + K \ \& \ p + K \leq a + s$
<proof>

lemma *word-succ-rbl*:

$$\text{to-bl } w = \text{bl} \implies \text{to-bl } (\text{word-succ } w) = (\text{rev } (\text{rbl-succ } (\text{rev } \text{bl})))$$

<proof>

lemma *word-pred-rbl*:

$$\text{to-bl } w = \text{bl} \implies \text{to-bl } (\text{word-pred } w) = (\text{rev } (\text{rbl-pred } (\text{rev } \text{bl})))$$

<proof>

lemma *word-add-rbl*:

$$\text{to-bl } v = \text{vbl} \implies \text{to-bl } w = \text{wbl} \implies$$

$$\text{to-bl } (v + w) = (\text{rev } (\text{rbl-add } (\text{rev } \text{vbl}) (\text{rev } \text{wbl})))$$

<proof>

lemma *word-mult-rbl*:

$$\text{to-bl } v = \text{vbl} \implies \text{to-bl } w = \text{wbl} \implies$$

$$\text{to-bl } (v * w) = (\text{rev } (\text{rbl-mult } (\text{rev } \text{vbl}) (\text{rev } \text{wbl})))$$

<proof>

lemma *rtb-rbl-ariths*:

$$\text{rev } (\text{to-bl } w) = \text{ys} \implies \text{rev } (\text{to-bl } (\text{word-succ } w)) = \text{rbl-succ } \text{ys}$$

$$\text{rev } (\text{to-bl } w) = \text{ys} \implies \text{rev } (\text{to-bl } (\text{word-pred } w)) = \text{rbl-pred } \text{ys}$$

$$\text{rev } (\text{to-bl } v) = \text{ys} \implies \text{rev } (\text{to-bl } w) = \text{xs} \implies \text{rev } (\text{to-bl } (v * w)) = \text{rbl-mult } \text{ys } \text{xs}$$

$$\text{rev } (\text{to-bl } v) = \text{ys} \implies \text{rev } (\text{to-bl } w) = \text{xs} \implies \text{rev } (\text{to-bl } (v + w)) = \text{rbl-add } \text{ys } \text{xs}$$

<proof>

16.20 Arithmetic type class instantiations

lemmas *word-le-0-iff* [*simp*] =

$$\text{word-zero-le } [\text{THEN } \text{leD}, \text{ THEN } \text{linorder-antisym-conv1}]$$

lemma *word-of-int-nat*:

$$0 \leq x \implies \text{word-of-int } x = \text{of-nat } (\text{nat } x)$$

<proof>

lemma *iszero-word-no* [*simp*]:

$$\text{iszero } (\text{numeral } \text{bin} :: 'a :: \text{len } \text{word}) =$$

$$\text{iszero } (\text{bintrunc } (\text{len-of } \text{TYPE}('a)) (\text{numeral } \text{bin}))$$

<proof>

Use *iszero* to simplify equalities between word numerals.

lemmas *word-eq-numeral-iff-iszero* [*simp*] =

$$\text{eq-numeral-iff-iszero } [\text{where } 'a = 'a :: \text{len } \text{word}]$$

16.21 Word and nat

lemma *td-ext-unat* [*OF refl*]:

$n = \text{len-of TYPE } ('a :: \text{len}) \implies$
 $\text{td-ext } (\text{unat} :: 'a \text{ word} \implies \text{nat}) \text{ of-nat}$
 $(\text{unats } n) (\%i. i \text{ mod } 2 \wedge n)$
 $\langle \text{proof} \rangle$

lemmas $\text{unat-of-nat} = \text{td-ext-unat} \text{ [THEN td-ext.eq-norm]}$

interpretation word-unat :
 $\text{td-ext } \text{unat} :: 'a :: \text{len word} \implies \text{nat}$
 of-nat
 $\text{unats } (\text{len-of TYPE } ('a :: \text{len}))$
 $\%i. i \text{ mod } 2 \wedge \text{len-of TYPE } ('a :: \text{len})$
 $\langle \text{proof} \rangle$

lemmas $\text{td-unat} = \text{word-unat.td-thm}$

lemmas $\text{unat-lt2p} \text{ [iff]} = \text{word-unat.Rep} \text{ [unfolded unats-def mem-Collect-eq]}$

lemma unat-le : $y \leq \text{unat } (z :: 'a :: \text{len word}) \implies y : \text{unats } (\text{len-of TYPE } ('a))$
 $\langle \text{proof} \rangle$

lemma word-nchotomy :
 $\text{ALL } w. \text{EX } n. (w :: 'a :: \text{len word}) = \text{of-nat } n \ \& \ n < 2 \wedge \text{len-of TYPE } ('a)$
 $\langle \text{proof} \rangle$

lemma of-nat-eq :
fixes $w :: 'a :: \text{len word}$
shows $(\text{of-nat } n = w) = (\exists q. n = \text{unat } w + q * 2 \wedge \text{len-of TYPE } ('a))$
 $\langle \text{proof} \rangle$

lemma of-nat-eq-size :
 $(\text{of-nat } n = w) = (\text{EX } q. n = \text{unat } w + q * 2 \wedge \text{size } w)$
 $\langle \text{proof} \rangle$

lemma of-nat-0 :
 $(\text{of-nat } m = (0 :: 'a :: \text{len word})) = (\exists q. m = q * 2 \wedge \text{len-of TYPE } ('a))$
 $\langle \text{proof} \rangle$

lemma $\text{of-nat-2p} \text{ [simp]}$:
 $\text{of-nat } (2 \wedge \text{len-of TYPE } ('a)) = (0 :: 'a :: \text{len word})$
 $\langle \text{proof} \rangle$

lemma of-nat-gt-0 : $\text{of-nat } k \sim = 0 \implies 0 < k$
 $\langle \text{proof} \rangle$

lemma of-nat-neq-0 :
 $0 < k \implies k < 2 \wedge \text{len-of TYPE } ('a :: \text{len}) \implies \text{of-nat } k \sim = (0 :: 'a \text{ word})$
 $\langle \text{proof} \rangle$

lemma *Abs-fnat-hom-add:*

$$\text{of-nat } a + \text{of-nat } b = \text{of-nat } (a + b)$$

<proof>

lemma *Abs-fnat-hom-mult:*

$$\text{of-nat } a * \text{of-nat } b = (\text{of-nat } (a * b) :: 'a :: \text{len word})$$

<proof>

lemma *Abs-fnat-hom-Suc:*

$$\text{word-succ } (\text{of-nat } a) = \text{of-nat } (\text{Suc } a)$$

<proof>

lemma *Abs-fnat-hom-0: (0::'a::len word) = of-nat 0*

<proof>

lemma *Abs-fnat-hom-1: (1::'a::len word) = of-nat (Suc 0)*

<proof>

lemmas *Abs-fnat-homs =*

Abs-fnat-hom-add Abs-fnat-hom-mult Abs-fnat-hom-Suc

Abs-fnat-hom-0 Abs-fnat-hom-1

lemma *word-arith-nat-add:*

$$a + b = \text{of-nat } (\text{unat } a + \text{unat } b)$$

<proof>

lemma *word-arith-nat-mult:*

$$a * b = \text{of-nat } (\text{unat } a * \text{unat } b)$$

<proof>

lemma *word-arith-nat-Suc:*

$$\text{word-succ } a = \text{of-nat } (\text{Suc } (\text{unat } a))$$

<proof>

lemma *word-arith-nat-div:*

$$a \text{ div } b = \text{of-nat } (\text{unat } a \text{ div } \text{unat } b)$$

<proof>

lemma *word-arith-nat-mod:*

$$a \text{ mod } b = \text{of-nat } (\text{unat } a \text{ mod } \text{unat } b)$$

<proof>

lemmas *word-arith-nat-defs =*

word-arith-nat-add word-arith-nat-mult

word-arith-nat-Suc Abs-fnat-hom-0

Abs-fnat-hom-1 word-arith-nat-div

word-arith-nat-mod

lemma *unat-cong: x = y \implies unat x = unat y*

<proof>

lemmas *unat-word-ariths* = *word-arith-nat-defs*
 [THEN trans [OF unat-cong unat-of-nat]]

lemmas *word-sub-less-iff* = *word-sub-le-iff*
 [unfolded linorder-not-less [symmetric] Not-eq-iff]

lemma *unat-add-lem*:
 (*unat x + unat y < 2 ^ len-of TYPE('a)*) =
 (*unat (x + y :: 'a :: len word) = unat x + unat y*)
<proof>

lemma *unat-mult-lem*:
 (*unat x * unat y < 2 ^ len-of TYPE('a)*) =
 (*unat (x * y :: 'a :: len word) = unat x * unat y*)
<proof>

lemmas *unat-plus-if'* = trans [OF unat-word-ariths(1) mod-nat-add, simplified]

lemma *le-no-overflow*:
x <= b ==> a <= a + b ==> x <= a + (b :: 'a :: len0 word)
<proof>

lemmas *un-ui-le* = trans [OF word-le-nat-alt [symmetric] word-le-def]

lemma *unat-sub-if-size*:
unat (x - y) = (if unat y <= unat x
then unat x - unat y
else unat x + 2 ^ size x - unat y)
<proof>

lemmas *unat-sub-if'* = *unat-sub-if-size* [unfolded word-size]

lemma *unat-div*: *unat ((x :: 'a :: len word) div y) = unat x div unat y*
<proof>

lemma *unat-mod*: *unat ((x :: 'a :: len word) mod y) = unat x mod unat y*
<proof>

lemma *uint-div*: *uint ((x :: 'a :: len word) div y) = uint x div uint y*
<proof>

lemma *uint-mod*: *uint ((x :: 'a :: len word) mod y) = uint x mod uint y*
<proof>

16.22 Definition of *unat-arith* tactic

lemma *unat-split*:

fixes $x :: 'a :: \text{len word}$
shows $P (\text{unat } x) =$
 $(\text{ALL } n. \text{ of-nat } n = x \ \& \ n < 2^{\text{len-of TYPE('a)}} \longrightarrow P \ n)$
 $\langle \text{proof} \rangle$

lemma *unat-split-asm*:
fixes $x :: 'a :: \text{len word}$
shows $P (\text{unat } x) =$
 $(\sim (\text{EX } n. \text{ of-nat } n = x \ \& \ n < 2^{\text{len-of TYPE('a)}} \ \& \ \sim P \ n))$
 $\langle \text{proof} \rangle$

lemmas *of-nat-inverse* =
 $\text{word-unat.Abs-inverse'}$ [*rotated, unfolded unats-def, simplified*]

lemmas *unat-splits* = *unat-split unat-split-asm*

lemmas *unat-arith-simps* =
word-le-nat-alt word-less-nat-alt
word-unat.Rep-inject [*symmetric*]
unat-sub-if' unat-plus-if' unat-div unat-mod

$\langle \text{ML} \rangle$

lemma *no-plus-overflow-unat-size*:
 $((x :: 'a :: \text{len word}) \leq x + y) = (\text{unat } x + \text{unat } y < 2^{\text{size } x})$
 $\langle \text{proof} \rangle$

lemmas *no-olen-add-nat* = *no-plus-overflow-unat-size* [*unfolded word-size*]

lemmas *unat-plus-simple* = *trans* [*OF no-olen-add-nat unat-add-lem*]

lemma *word-div-mult*:
 $(0 :: 'a :: \text{len word}) < y \implies \text{unat } x * \text{unat } y < 2^{\text{len-of TYPE('a)}} \implies$
 $x * y \text{ div } y = x$
 $\langle \text{proof} \rangle$

lemma *div-lt'*: $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies$
 $\text{unat } i * \text{unat } x < 2^{\text{len-of TYPE('a)}}$
 $\langle \text{proof} \rangle$

lemmas *div-lt''* = *order-less-imp-le* [*THEN div-lt'*]

lemma *div-lt-mult*: $(i :: 'a :: \text{len word}) < k \text{ div } x \implies 0 < x \implies i * x < k$
 $\langle \text{proof} \rangle$

lemma *div-le-mult*:
 $(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies 0 < x \implies i * x \leq k$
 $\langle \text{proof} \rangle$

lemma *div-lt-uint'*:

$(i :: 'a :: \text{len word}) \leq k \text{ div } x \implies \text{uint } i * \text{uint } x < 2 \wedge \text{len-of TYPE}('a)$
 ⟨proof⟩

lemmas *div-lt-uint''* = *order-less-imp-le* [THEN *div-lt-uint'*]

lemma *word-le-exists'*:

$(x :: 'a :: \text{len0 word}) \leq y \implies$
 $(\text{EX } z. y = x + z \ \& \ \text{uint } x + \text{uint } z < 2 \wedge \text{len-of TYPE}('a))$
 ⟨proof⟩

lemmas *plus-minus-not-NULL* = *order-less-imp-le* [THEN *plus-minus-not-NULL-ab*]

lemmas *plus-minus-no-overflow* =
order-less-imp-le [THEN *plus-minus-no-overflow-ab*]

lemmas *mcs* = *word-less-minus-cancel* *word-less-minus-mono-left*
word-le-minus-cancel *word-le-minus-mono-left*

lemmas *word-l-diffs* = *mcs* [where $y = w + x$, *unfolded add-diff-cancel*] **for** $w \ x$

lemmas *word-diff-ls* = *mcs* [where $z = w + x$, *unfolded add-diff-cancel*] **for** $w \ x$

lemmas *word-plus-mcs* = *word-diff-ls* [where $y = v + x$, *unfolded add-diff-cancel*]
for $v \ x$

lemmas *le-unat-uo* = *unat-le* [THEN *word-unat.Abs-inverse*]

lemmas *thd* = *refl* [THEN [2] *split-div-lemma* [THEN *iffD2*], THEN *conjunct1*]

lemmas *uno-simps* [THEN *le-unat-uo*] = *mod-le-divisor* *div-le-dividend* *dtle*

lemma *word-mod-div-equality*:

$(n \text{ div } b) * b + (n \text{ mod } b) = (n :: 'a :: \text{len word})$
 ⟨proof⟩

lemma *word-div-mult-le*: $a \text{ div } b * b \leq (a :: 'a :: \text{len word})$

⟨proof⟩

lemma *word-mod-less-divisor*: $0 < n \implies m \text{ mod } n < (n :: 'a :: \text{len word})$

⟨proof⟩

lemma *word-of-int-power-hom*:

$\text{word-of-int } a \wedge n = (\text{word-of-int } (a \wedge n) :: 'a :: \text{len word})$
 ⟨proof⟩

lemma *word-arith-power-alt*:

$a \wedge n = (\text{word-of-int } (\text{uint } a \wedge n) :: 'a :: \text{len word})$
 ⟨proof⟩

lemma *of-bl-length-less*:

$length\ x = k \implies k < len\text{-of}\ TYPE('a) \implies (of\text{-bl}\ x :: 'a :: len\ word) < 2 \wedge k$
 ⟨*proof*⟩

16.23 Cardinality, finiteness of set of words

instance *word* :: (*len0*) *finite*

⟨*proof*⟩

lemma *card-word*: $CARD('a::len0\ word) = 2 \wedge len\text{-of}\ TYPE('a)$

⟨*proof*⟩

lemma *card-word-size*:

$card\ (UNIV :: 'a :: len0\ word\ set) = (2 \wedge size\ (x :: 'a\ word))$
 ⟨*proof*⟩

16.24 Bitwise Operations on Words

lemmas *bin-log-bintrs* = *bin-trunc-not bin-trunc-xor bin-trunc-and bin-trunc-or*

lemmas *wils1* = *bin-log-bintrs* [*THEN word-ubin.norm-eq-iff* [*THEN iffD1*],
folded word-ubin.eq-norm, THEN eq-reflection]

lemmas *word-log-binary-defs* =

word-and-def word-or-def word-xor-def

lemma *word-wi-log-defs*:

NOT word-of-int a = *word-of-int (NOT a)*
word-of-int a AND word-of-int b = *word-of-int (a AND b)*
word-of-int a OR word-of-int b = *word-of-int (a OR b)*
word-of-int a XOR word-of-int b = *word-of-int (a XOR b)*
 ⟨*proof*⟩

lemma *word-no-log-defs* [*simp*]:

NOT (numeral a) = *word-of-int (NOT (numeral a))*
NOT (– numeral a) = *word-of-int (NOT (– numeral a))*
numeral a AND numeral b = *word-of-int (numeral a AND numeral b)*
numeral a AND – numeral b = *word-of-int (numeral a AND – numeral b)*
– numeral a AND numeral b = *word-of-int (– numeral a AND numeral b)*
– numeral a AND – numeral b = *word-of-int (– numeral a AND – numeral b)*
numeral a OR numeral b = *word-of-int (numeral a OR numeral b)*
numeral a OR – numeral b = *word-of-int (numeral a OR – numeral b)*
– numeral a OR numeral b = *word-of-int (– numeral a OR numeral b)*
– numeral a OR – numeral b = *word-of-int (– numeral a OR – numeral b)*
numeral a XOR numeral b = *word-of-int (numeral a XOR numeral b)*

$\text{numeral } a \text{ XOR } - \text{ numeral } b = \text{word-of-int } (\text{numeral } a \text{ XOR } - \text{ numeral } b)$
 $- \text{ numeral } a \text{ XOR numeral } b = \text{word-of-int } (- \text{ numeral } a \text{ XOR numeral } b)$
 $- \text{ numeral } a \text{ XOR } - \text{ numeral } b = \text{word-of-int } (- \text{ numeral } a \text{ XOR } - \text{ numeral } b)$
 ⟨proof⟩

Special cases for when one of the arguments equals 1.

lemma *word-bitwise-1-simps* [simp]:

$\text{NOT } (1::'a::\text{len0 word}) = -2$
 $1 \text{ AND numeral } b = \text{word-of-int } (1 \text{ AND numeral } b)$
 $1 \text{ AND } - \text{ numeral } b = \text{word-of-int } (1 \text{ AND } - \text{ numeral } b)$
 $\text{numeral } a \text{ AND } 1 = \text{word-of-int } (\text{numeral } a \text{ AND } 1)$
 $- \text{ numeral } a \text{ AND } 1 = \text{word-of-int } (- \text{ numeral } a \text{ AND } 1)$
 $1 \text{ OR numeral } b = \text{word-of-int } (1 \text{ OR numeral } b)$
 $1 \text{ OR } - \text{ numeral } b = \text{word-of-int } (1 \text{ OR } - \text{ numeral } b)$
 $\text{numeral } a \text{ OR } 1 = \text{word-of-int } (\text{numeral } a \text{ OR } 1)$
 $- \text{ numeral } a \text{ OR } 1 = \text{word-of-int } (- \text{ numeral } a \text{ OR } 1)$
 $1 \text{ XOR numeral } b = \text{word-of-int } (1 \text{ XOR numeral } b)$
 $1 \text{ XOR } - \text{ numeral } b = \text{word-of-int } (1 \text{ XOR } - \text{ numeral } b)$
 $\text{numeral } a \text{ XOR } 1 = \text{word-of-int } (\text{numeral } a \text{ XOR } 1)$
 $- \text{ numeral } a \text{ XOR } 1 = \text{word-of-int } (- \text{ numeral } a \text{ XOR } 1)$
 ⟨proof⟩

Special cases for when one of the arguments equals -1.

lemma *word-bitwise-m1-simps* [simp]:

$\text{NOT } (-1::'a::\text{len0 word}) = 0$
 $(-1::'a::\text{len0 word}) \text{ AND } x = x$
 $x \text{ AND } (-1::'a::\text{len0 word}) = x$
 $(-1::'a::\text{len0 word}) \text{ OR } x = -1$
 $x \text{ OR } (-1::'a::\text{len0 word}) = -1$
 $(-1::'a::\text{len0 word}) \text{ XOR } x = \text{NOT } x$
 $x \text{ XOR } (-1::'a::\text{len0 word}) = \text{NOT } x$
 ⟨proof⟩

lemma *uint-or*: $\text{uint } (x \text{ OR } y) = (\text{uint } x) \text{ OR } (\text{uint } y)$

⟨proof⟩

lemma *uint-and*: $\text{uint } (x \text{ AND } y) = (\text{uint } x) \text{ AND } (\text{uint } y)$

⟨proof⟩

lemma *test-bit-wi* [simp]:

$(\text{word-of-int } x::'a::\text{len0 word}) !! n \longleftrightarrow n < \text{len-of TYPE('a)} \wedge \text{bin-nth } x n$
 ⟨proof⟩

lemma *word-test-bit-transfer* [transfer-rule]:

$(\text{rel-fun pcr-word } (\text{rel-fun } op = op =))$
 $(\lambda x n. n < \text{len-of TYPE('a)} \wedge \text{bin-nth } x n) (\text{test-bit } :: 'a::\text{len0 word} \Rightarrow -)$
 ⟨proof⟩

lemma *word-ops-nth-size*:

$n < \text{size } (x :: 'a :: \text{len0 word}) \implies$
 $(x \text{ OR } y) !! n = (x !! n | y !! n) \&$
 $(x \text{ AND } y) !! n = (x !! n \& y !! n) \&$
 $(x \text{ XOR } y) !! n = (x !! n \sim y !! n) \&$
 $(\text{NOT } x) !! n = (\sim x !! n)$
 ⟨proof⟩

lemma *word-ao-nth*:

fixes $x :: 'a :: \text{len0 word}$
shows $(x \text{ OR } y) !! n = (x !! n | y !! n) \&$
 $(x \text{ AND } y) !! n = (x !! n \& y !! n)$
 ⟨proof⟩

lemma *test-bit-numeral* [simp]:

$(\text{numeral } w :: 'a :: \text{len0 word}) !! n \longleftrightarrow$
 $n < \text{len-of TYPE('a)} \wedge \text{bin-nth } (\text{numeral } w) n$
 ⟨proof⟩

lemma *test-bit-neg-numeral* [simp]:

$(- \text{numeral } w :: 'a :: \text{len0 word}) !! n \longleftrightarrow$
 $n < \text{len-of TYPE('a)} \wedge \text{bin-nth } (- \text{numeral } w) n$
 ⟨proof⟩

lemma *test-bit-1* [simp]: $(1 :: 'a :: \text{len word}) !! n \longleftrightarrow n = 0$

⟨proof⟩

lemma *nth-0* [simp]: $\sim (0 :: 'a :: \text{len0 word}) !! n$

⟨proof⟩

lemma *nth-minus1* [simp]: $(-1 :: 'a :: \text{len0 word}) !! n \longleftrightarrow n < \text{len-of TYPE('a)}$

⟨proof⟩

lemmas *bwsimps* =

wi-hom-add
word-wi-log-defs

lemma *word-bw-assocs*:

fixes $x :: 'a :: \text{len0 word}$
shows
 $(x \text{ AND } y) \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $(x \text{ OR } y) \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $(x \text{ XOR } y) \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
 ⟨proof⟩

lemma *word-bw-comms*:

fixes $x :: 'a :: \text{len0 word}$
shows

$x \text{ AND } y = y \text{ AND } x$
 $x \text{ OR } y = y \text{ OR } x$
 $x \text{ XOR } y = y \text{ XOR } x$
 ⟨proof⟩

lemma *word-bw-lcs*:

fixes $x :: 'a::len0 \text{ word}$
shows
 $y \text{ AND } x \text{ AND } z = x \text{ AND } y \text{ AND } z$
 $y \text{ OR } x \text{ OR } z = x \text{ OR } y \text{ OR } z$
 $y \text{ XOR } x \text{ XOR } z = x \text{ XOR } y \text{ XOR } z$
 ⟨proof⟩

lemma *word-log-esimps* [simp]:

fixes $x :: 'a::len0 \text{ word}$
shows
 $x \text{ AND } 0 = 0$
 $x \text{ AND } -1 = x$
 $x \text{ OR } 0 = x$
 $x \text{ OR } -1 = -1$
 $x \text{ XOR } 0 = x$
 $x \text{ XOR } -1 = \text{NOT } x$
 $0 \text{ AND } x = 0$
 $-1 \text{ AND } x = x$
 $0 \text{ OR } x = x$
 $-1 \text{ OR } x = -1$
 $0 \text{ XOR } x = x$
 $-1 \text{ XOR } x = \text{NOT } x$
 ⟨proof⟩

lemma *word-not-dist*:

fixes $x :: 'a::len0 \text{ word}$
shows
 $\text{NOT } (x \text{ OR } y) = \text{NOT } x \text{ AND } \text{NOT } y$
 $\text{NOT } (x \text{ AND } y) = \text{NOT } x \text{ OR } \text{NOT } y$
 ⟨proof⟩

lemma *word-bw-same*:

fixes $x :: 'a::len0 \text{ word}$
shows
 $x \text{ AND } x = x$
 $x \text{ OR } x = x$
 $x \text{ XOR } x = 0$
 ⟨proof⟩

lemma *word-ao-absorbs* [simp]:

fixes $x :: 'a::len0 \text{ word}$
shows
 $x \text{ AND } (y \text{ OR } x) = x$

$x \text{ OR } y \text{ AND } x = x$
 $x \text{ AND } (x \text{ OR } y) = x$
 $y \text{ AND } x \text{ OR } x = x$
 $(y \text{ OR } x) \text{ AND } x = x$
 $x \text{ OR } x \text{ AND } y = x$
 $(x \text{ OR } y) \text{ AND } x = x$
 $x \text{ AND } y \text{ OR } x = x$
 ⟨proof⟩

lemma *word-not-not* [simp]:
 $\text{NOT NOT } (x :: 'a :: \text{len0 word}) = x$
 ⟨proof⟩

lemma *word-ao-dist*:
fixes $x :: 'a :: \text{len0 word}$
shows $(x \text{ OR } y) \text{ AND } z = x \text{ AND } z \text{ OR } y \text{ AND } z$
 ⟨proof⟩

lemma *word-oa-dist*:
fixes $x :: 'a :: \text{len0 word}$
shows $x \text{ AND } y \text{ OR } z = (x \text{ OR } z) \text{ AND } (y \text{ OR } z)$
 ⟨proof⟩

lemma *word-add-not* [simp]:
fixes $x :: 'a :: \text{len0 word}$
shows $x + \text{NOT } x = -1$
 ⟨proof⟩

lemma *word-plus-and-or* [simp]:
fixes $x :: 'a :: \text{len0 word}$
shows $(x \text{ AND } y) + (x \text{ OR } y) = x + y$
 ⟨proof⟩

lemma *lea*:
fixes $x :: 'a :: \text{len0 word}$
shows $(w = (x \text{ OR } y)) \implies (y = (w \text{ AND } y))$ ⟨proof⟩

lemma *leao*:
fixes $x' :: 'a :: \text{len0 word}$
shows $(w' = (x' \text{ AND } y')) \implies (x' = (x' \text{ OR } w'))$ ⟨proof⟩

lemma *word-ao-equiv*:
fixes $w w' :: 'a :: \text{len0 word}$
shows $(w = w \text{ OR } w') = (w' = w \text{ AND } w')$
 ⟨proof⟩

lemma *le-word-or2*: $x \leq x \text{ OR } (y :: 'a :: \text{len0 word})$
 ⟨proof⟩

lemmas *le-word-or1* = *xtr3* [OF *word-bw-comms* (2) *le-word-or2*]

lemmas *word-and-le1* = *xtr3* [*OF word-ao-absorbs* (4) [*symmetric*] *le-word-or2*]
lemmas *word-and-le2* = *xtr3* [*OF word-ao-absorbs* (8) [*symmetric*] *le-word-or2*]

lemma *bl-word-not*: $to\text{-}bl\ (NOT\ w) = map2\ Not\ (to\text{-}bl\ w)$
 ⟨*proof*⟩

lemma *bl-word-xor*: $to\text{-}bl\ (v\ XOR\ w) = map2\ op\ \sim = (to\text{-}bl\ v)\ (to\text{-}bl\ w)$
 ⟨*proof*⟩

lemma *bl-word-or*: $to\text{-}bl\ (v\ OR\ w) = map2\ op\ |\ (to\text{-}bl\ v)\ (to\text{-}bl\ w)$
 ⟨*proof*⟩

lemma *bl-word-and*: $to\text{-}bl\ (v\ AND\ w) = map2\ op\ \&\ (to\text{-}bl\ v)\ (to\text{-}bl\ w)$
 ⟨*proof*⟩

lemma *word-lsb-alt*: $lsb\ (w::'a::len0\ word) = test\text{-}bit\ w\ 0$
 ⟨*proof*⟩

lemma *word-lsb-1-0* [*simp*]: $lsb\ (1::'a::len\ word) \&\ \sim\ lsb\ (0::'b::len0\ word)$
 ⟨*proof*⟩

lemma *word-lsb-last*: $lsb\ (w::'a::len\ word) = last\ (to\text{-}bl\ w)$
 ⟨*proof*⟩

lemma *word-lsb-int*: $lsb\ w = (uint\ w\ mod\ 2 = 1)$
 ⟨*proof*⟩

lemma *word-msb-sint*: $msb\ w = (sint\ w < 0)$
 ⟨*proof*⟩

lemma *msb-word-of-int*:
 $msb\ (word\text{-}of\text{-}int\ x::'a::len\ word) = bin\text{-}nth\ x\ (len\text{-}of\ TYPE('a) - 1)$
 ⟨*proof*⟩

lemma *word-msb-numeral* [*simp*]:
 $msb\ (numeral\ w::'a::len\ word) = bin\text{-}nth\ (numeral\ w)\ (len\text{-}of\ TYPE('a) - 1)$
 ⟨*proof*⟩

lemma *word-msb-neg-numeral* [*simp*]:
 $msb\ (-\ numeral\ w::'a::len\ word) = bin\text{-}nth\ (-\ numeral\ w)\ (len\text{-}of\ TYPE('a) - 1)$
 ⟨*proof*⟩

lemma *word-msb-0* [*simp*]: $\neg\ msb\ (0::'a::len\ word)$
 ⟨*proof*⟩

lemma *word-msb-1* [*simp*]: $msb\ (1::'a::len\ word) \longleftrightarrow len\text{-}of\ TYPE('a) = 1$
 ⟨*proof*⟩

lemma *word-msb-nth*:

$msb (w :: 'a :: len \text{ word}) = bin\text{-}nth (uint\ w) (len\text{-}of\ TYPE('a) - 1)$
 $\langle proof \rangle$

lemma *word-msb-alt*: $msb (w :: 'a :: len \text{ word}) = hd (to\text{-}bl\ w)$

$\langle proof \rangle$

lemma *word-set-nth* [*simp*]:

$set\text{-}bit\ w\ n (test\text{-}bit\ w\ n) = (w :: 'a :: len0 \text{ word})$
 $\langle proof \rangle$

lemma *bin-nth-uint'*:

$bin\text{-}nth (uint\ w)\ n = (rev (bin\text{-}to\text{-}bl (size\ w) (uint\ w)) ! n \ \&\ n < size\ w)$
 $\langle proof \rangle$

lemmas *bin-nth-uint = bin-nth-uint'* [*unfolded word-size*]

lemma *test-bit-bl*: $w !! n = (rev (to\text{-}bl\ w) ! n \ \&\ n < size\ w)$

$\langle proof \rangle$

lemma *to-bl-nth*: $n < size\ w \implies to\text{-}bl\ w ! n = w !! (size\ w - Suc\ n)$

$\langle proof \rangle$

lemma *test-bit-set*:

fixes $w :: 'a :: len0 \text{ word}$
shows $(set\text{-}bit\ w\ n\ x) !! n = (n < size\ w \ \&\ x)$
 $\langle proof \rangle$

lemma *test-bit-set-gen*:

fixes $w :: 'a :: len0 \text{ word}$
shows $test\text{-}bit (set\text{-}bit\ w\ n\ x)\ m =$
 $(if\ m = n\ then\ n < size\ w \ \&\ x\ else\ test\text{-}bit\ w\ m)$
 $\langle proof \rangle$

lemma *of-bl-rep-False*: $of\text{-}bl (replicate\ n\ False\ @\ bs) = of\text{-}bl\ bs$

$\langle proof \rangle$

lemma *msb-nth*:

fixes $w :: 'a :: len \text{ word}$
shows $msb\ w = w !! (len\text{-}of\ TYPE('a) - 1)$
 $\langle proof \rangle$

lemmas $msb0 = len\text{-}gt\ 0$ [*THEN diff-Suc-less, THEN word-ops-nth-size* [*unfolded word-size*]]

lemmas $msb1 = msb0$ [**where** $i = 0$]

lemmas $word\text{-}ops\text{-}msb = msb1$ [*unfolded msb-nth* [*symmetric, unfolded One-nat-def*]]

lemmas $lsb0 = len\text{-}gt\ 0$ [*THEN word-ops-nth-size* [*unfolded word-size*]]

lemmas $word\text{-}ops\text{-}lsb = lsb0$ [*unfolded word-lsb-alt*]

lemma *td-ext-nth* [*OF refl refl refl, unfolded word-size*]:
 $n = \text{size } (w :: 'a :: \text{len0 word}) \implies \text{ofn} = \text{set-bits} \implies [w, \text{ofn } g] = l \implies$
 $\text{td-ext test-bit ofn } \{f. \text{ALL } i. f \ i \ \longrightarrow \ i < n\} \ (\%h \ i. \ h \ i \ \& \ i < n)$
<proof>

interpretation *test-bit*:
 $\text{td-ext op !!} :: 'a :: \text{len0 word} \Rightarrow \text{nat} \Rightarrow \text{bool}$
 set-bits
 $\{f. \forall i. f \ i \ \longrightarrow \ i < \text{len-of TYPE}('a :: \text{len0})\}$
 $(\lambda h \ i. \ h \ i \ \wedge \ i < \text{len-of TYPE}('a :: \text{len0}))$
<proof>

lemmas $\text{td-nth} = \text{test-bit.td-thm}$

lemma *word-set-set-same* [*simp*]:
fixes $w :: 'a :: \text{len0 word}$
shows $\text{set-bit } (\text{set-bit } w \ n \ x) \ n \ y = \text{set-bit } w \ n \ y$
<proof>

lemma *word-set-set-diff*:
fixes $w :: 'a :: \text{len0 word}$
assumes $m \ \sim = \ n$
shows $\text{set-bit } (\text{set-bit } w \ m \ x) \ n \ y = \text{set-bit } (\text{set-bit } w \ n \ y) \ m \ x$
<proof>

lemma *nth-sint*:
fixes $w :: 'a :: \text{len word}$
defines $l \equiv \text{len-of TYPE } ('a)$
shows $\text{bin-nth } (\text{sint } w) \ n = (\text{if } n < l - 1 \text{ then } w \ !! \ n \ \text{else } w \ !! \ (l - 1))$
<proof>

lemma *word-lsb-numeral* [*simp*]:
 $\text{lsb } (\text{numeral bin} :: 'a :: \text{len word}) \ \longleftrightarrow \ \text{bin-last } (\text{numeral bin})$
<proof>

lemma *word-lsb-neg-numeral* [*simp*]:
 $\text{lsb } (- \text{numeral bin} :: 'a :: \text{len word}) \ \longleftrightarrow \ \text{bin-last } (- \text{numeral bin})$
<proof>

lemma *set-bit-word-of-int*:
 $\text{set-bit } (\text{word-of-int } x) \ n \ b = \text{word-of-int } (\text{bin-sc } n \ b \ x)$
<proof>

lemma *word-set-numeral* [*simp*]:
 $\text{set-bit } (\text{numeral bin} :: 'a :: \text{len0 word}) \ n \ b =$
 $\text{word-of-int } (\text{bin-sc } n \ b \ (\text{numeral bin}))$
<proof>

lemma *word-set-neg-numeral* [simp]:

set-bit ($-$ numeral bin::*a*::len0 word) *n b* =
word-of-int (bin-sc *n b* ($-$ numeral bin))
 ⟨proof⟩

lemma *word-set-bit-0* [simp]:

set-bit 0 *n b* = *word-of-int* (bin-sc *n b* 0)
 ⟨proof⟩

lemma *word-set-bit-1* [simp]:

set-bit 1 *n b* = *word-of-int* (bin-sc *n b* 1)
 ⟨proof⟩

lemma *setBit-no* [simp]:

setBit (numeral bin) *n* = *word-of-int* (bin-sc *n* True (numeral bin))
 ⟨proof⟩

lemma *clearBit-no* [simp]:

clearBit (numeral bin) *n* = *word-of-int* (bin-sc *n* False (numeral bin))
 ⟨proof⟩

lemma *to-bl-n1*:

to-bl (-1 ::*a*::len0 word) = *replicate* (len-of TYPE (*a*)) True
 ⟨proof⟩

lemma *word-msb-n1* [simp]: *msb* (-1 ::*a*::len word)

⟨proof⟩

lemma *word-set-nth-iff*:

(*set-bit* *w n b* = *w*) = (*w* !! *n* = *b* | *n* >= size (*w*::*a*::len0 word))
 ⟨proof⟩

lemma *test-bit-2p*:

(*word-of-int* (2 ^ *n*)::*a*::len word) !! *m* \longleftrightarrow *m* = *n* \wedge *m* < len-of TYPE(*a*)
 ⟨proof⟩

lemma *nth-w2p*:

((2 ::*a*::len word) ^ *n*) !! *m* \longleftrightarrow *m* = *n* \wedge *m* < len-of TYPE(*a*::len)
 ⟨proof⟩

lemma *uint-2p*:

(0 ::*a*::len word) < 2 ^ *n* \implies *uint* (2 ^ *n*::*a*::len word) = 2 ^ *n*
 ⟨proof⟩

lemma *word-of-int-2p*: (*word-of-int* (2 ^ *n*)::*a*::len word) = 2 ^ *n*

⟨proof⟩

lemma *bang-is-le*: *x* !! *m* \implies 2 ^ *m* <= (*x* :: *a* :: len word)

⟨proof⟩

lemma *word-clr-le*:
fixes $w :: 'a::len0\ word$
shows $w \geq \text{set-bit } w\ n\ \text{False}$
 $\langle \text{proof} \rangle$

lemma *word-set-ge*:
fixes $w :: 'a::len\ word$
shows $w \leq \text{set-bit } w\ n\ \text{True}$
 $\langle \text{proof} \rangle$

16.25 Shifting, Rotating, and Splitting Words

lemma *shiffl1-wi* [simp]: $\text{shiffl1 } (\text{word-of-int } w) = \text{word-of-int } (w\ \text{BIT}\ \text{False})$
 $\langle \text{proof} \rangle$

lemma *shiffl1-numeral* [simp]:
 $\text{shiffl1 } (\text{numeral } w) = \text{numeral } (\text{Num.Bit0 } w)$
 $\langle \text{proof} \rangle$

lemma *shiffl1-neg-numeral* [simp]:
 $\text{shiffl1 } (-\ \text{numeral } w) = -\ \text{numeral } (\text{Num.Bit0 } w)$
 $\langle \text{proof} \rangle$

lemma *shiffl1-0* [simp]: $\text{shiffl1 } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *shiffl1-def-u*: $\text{shiffl1 } w = \text{word-of-int } (\text{uint } w\ \text{BIT}\ \text{False})$
 $\langle \text{proof} \rangle$

lemma *shiffl1-def-s*: $\text{shiffl1 } w = \text{word-of-int } (\text{sint } w\ \text{BIT}\ \text{False})$
 $\langle \text{proof} \rangle$

lemma *shiftr1-0* [simp]: $\text{shiftr1 } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *sshiftr1-0* [simp]: $\text{sshiftr1 } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *sshiftr1-n1* [simp]: $\text{sshiftr1 } (-\ 1) = -\ 1$
 $\langle \text{proof} \rangle$

lemma *shifl-0* [simp]: $(0::'a::len0\ word) \ll n = 0$
 $\langle \text{proof} \rangle$

lemma *shiftr-0* [simp]: $(0::'a::len0\ word) \gg n = 0$
 $\langle \text{proof} \rangle$

lemma *sshiftr-0* [simp]: $0 \ggg n = 0$

<proof>

lemma *sshiftr-n1* [*simp*] : $-1 \gg \gg n = -1$

<proof>

lemma *nth-shiftl1*: $\text{shiftl1 } w \text{ !! } n = (n < \text{size } w \ \& \ n > 0 \ \& \ w \text{ !! } (n - 1))$

<proof>

lemma *nth-shiftl'* [*rule-format*]:

ALL $n. ((w::'a::\text{len0 } \text{word}) << m) \text{ !! } n = (n < \text{size } w \ \& \ n \geq m \ \& \ w \text{ !! } (n - m))$

<proof>

lemmas *nth-shiftl = nth-shiftl'* [*unfolded word-size*]

lemma *nth-shiftr1*: $\text{shiftr1 } w \text{ !! } n = w \text{ !! } \text{Suc } n$

<proof>

lemma *nth-shiftr*:

$\bigwedge n. ((w::'a::\text{len0 } \text{word}) \gg m) \text{ !! } n = w \text{ !! } (n + m)$

<proof>

lemma *uint-shiftr1*: $\text{uint } (\text{shiftr1 } w) = \text{bin-rest } (\text{uint } w)$

<proof>

lemma *nth-sshiftr1*:

$\text{sshiftr1 } w \text{ !! } n = (\text{if } n = \text{size } w - 1 \text{ then } w \text{ !! } n \text{ else } w \text{ !! } \text{Suc } n)$

<proof>

lemma *nth-sshiftr* [*rule-format*] :

ALL $n. \text{sshiftr } w \ m \ \text{!! } n = (n < \text{size } w \ \&$

$(\text{if } n + m \geq \text{size } w \text{ then } w \ \text{!! } (\text{size } w - 1) \text{ else } w \ \text{!! } (n + m)))$

<proof>

lemma *shiftr1-div-2*: $\text{uint } (\text{shiftr1 } w) = \text{uint } w \ \text{div } 2$

<proof>

lemma *sshiftr1-div-2*: $\text{sint } (\text{sshiftr1 } w) = \text{sint } w \ \text{div } 2$

<proof>

lemma *shiftr-div-2n*: $\text{uint } (\text{shiftr } w \ n) = \text{uint } w \ \text{div } 2 \ ^n$

<proof>

lemma *sshiftr-div-2n*: $\text{sint } (\text{sshiftr } w \ n) = \text{sint } w \ \text{div } 2 \ ^n$

<proof>

16.25.1 shift functions in terms of lists of bools

lemmas *bshiftr1-numeral* [*simp*] =
bshiftr1-def [**where** *w=numeral w*, *unfolded to-bl-numeral*] **for** *w*

lemma *bshiftr1-bl*: $to-bl (bshiftr1\ b\ w) = b \# butlast (to-bl\ w)$
 ⟨*proof*⟩

lemma *shiftrl1-of-bl*: $shiftrl1 (of-bl\ bl) = of-bl (bl\ @\ [False])$
 ⟨*proof*⟩

lemma *shiftrl1-bl*: $shiftrl1 (w::'a::len0\ word) = of-bl (to-bl\ w\ @\ [False])$
 ⟨*proof*⟩

lemma *bl-shiftrl1*:
 $to-bl (shiftrl1 (w :: 'a :: len\ word)) = tl (to-bl\ w) @ [False]$
 ⟨*proof*⟩

lemma *bl-shiftrl1'*:
 $to-bl (shiftrl1\ w) = tl (to-bl\ w @ [False])$
 ⟨*proof*⟩

lemma *shiftr1-bl*: $shiftr1\ w = of-bl (butlast (to-bl\ w))$
 ⟨*proof*⟩

lemma *bl-shiftr1*:
 $to-bl (shiftr1 (w :: 'a :: len\ word)) = False \# butlast (to-bl\ w)$
 ⟨*proof*⟩

lemma *bl-shiftr1'*:
 $to-bl (shiftr1\ w) = butlast (False \# to-bl\ w)$
 ⟨*proof*⟩

lemma *shiftrl1-rev*:
 $shiftrl1\ w = word-reverse (shiftr1 (word-reverse\ w))$
 ⟨*proof*⟩

lemma *shiftrl-rev*:
 $shiftrl\ w\ n = word-reverse (shiftr (word-reverse\ w)\ n)$
 ⟨*proof*⟩

lemma *rev-shiftrl*: $word-reverse\ w \ll n = word-reverse (w \gg n)$
 ⟨*proof*⟩

lemma *shiftr-rev*: $w \gg n = word-reverse (word-reverse\ w \ll n)$
 ⟨*proof*⟩

lemma *rev-shiftr*: $word-reverse\ w \gg n = word-reverse (w \ll n)$

<proof>

lemma *bl-sshiftr1*:

$to-bl (sshiftr1 (w :: 'a :: len word)) = hd (to-bl w) \# butlast (to-bl w)$
<proof>

lemma *drop-shiftr*:

$drop n (to-bl ((w :: 'a :: len word) >> n)) = take (size w - n) (to-bl w)$
<proof>

lemma *drop-sshiftr*:

$drop n (to-bl ((w :: 'a :: len word) >>> n)) = take (size w - n) (to-bl w)$
<proof>

lemma *take-shiftr*:

$n \leq size w \implies take n (to-bl (w >> n)) = replicate n False$
<proof>

lemma *take-sshiftr'* [*rule-format*] :

$n \leq size (w :: 'a :: len word) \implies hd (to-bl (w >>> n)) = hd (to-bl w) \&$
 $take n (to-bl (w >>> n)) = replicate n (hd (to-bl w))$
<proof>

lemmas *hd-sshiftr = take-sshiftr'* [*THEN conjunct1*]

lemmas *take-sshiftr = take-sshiftr'* [*THEN conjunct2*]

lemma *atd-lem*: $take n xs = t \implies drop n xs = d \implies xs = t @ d$

<proof>

lemmas *bl-shiftr = atd-lem* [*OF take-shiftr drop-shiftr*]

lemmas *bl-sshiftr = atd-lem* [*OF take-sshiftr drop-sshiftr*]

lemma *shiftr-of-bl*: $of-bl bl \ll n = of-bl (bl @ replicate n False)$

<proof>

lemma *shiftr-bl*:

$(w :: 'a :: len0 word) \ll (n :: nat) = of-bl (to-bl w @ replicate n False)$
<proof>

lemmas *shiftr-numeral* [*simp*] = *shiftr-def* [**where** $w = numeral w$] **for** w

lemma *bl-shiftr*:

$to-bl (w \ll n) = drop n (to-bl w) @ replicate (min (size w) n) False$
<proof>

lemma *shiftr-zero-size*:

fixes $x :: 'a :: len0 word$

shows $size x \leq n \implies x \ll n = 0$

<proof>

lemma *shiffl1-2t*: $\text{shiffl1 } (w :: 'a :: \text{len word}) = 2 * w$
 ⟨proof⟩

lemma *shiffl1-p*: $\text{shiffl1 } (w :: 'a :: \text{len word}) = w + w$
 ⟨proof⟩

lemma *shiffl-t2n*: $\text{shiffl } (w :: 'a :: \text{len word}) \ n = 2 ^ n * w$
 ⟨proof⟩

lemma *shiftr1-bintr* [simp]:
 $(\text{shiftr1 } (\text{numeral } w) :: 'a :: \text{len0 word}) =$
 $\text{word-of-int } (\text{bin-rest } (\text{bintrunc } (\text{len-of TYPE } ('a)) (\text{numeral } w)))$
 ⟨proof⟩

lemma *sshiftr1-sbintr* [simp]:
 $(\text{sshiftr1 } (\text{numeral } w) :: 'a :: \text{len word}) =$
 $\text{word-of-int } (\text{bin-rest } (\text{sbintrunc } (\text{len-of TYPE } ('a) - 1) (\text{numeral } w)))$
 ⟨proof⟩

lemma *shiftr-no* [simp]:

$(\text{numeral } w :: 'a :: \text{len0 word}) \gg n = \text{word-of-int}$
 $((\text{bin-rest } ^n) (\text{bintrunc } (\text{len-of TYPE } ('a)) (\text{numeral } w)))$
 ⟨proof⟩

lemma *sshiftr-no* [simp]:

$(\text{numeral } w :: 'a :: \text{len word}) \ggg n = \text{word-of-int}$
 $((\text{bin-rest } ^n) (\text{sbintrunc } (\text{len-of TYPE } ('a) - 1) (\text{numeral } w)))$
 ⟨proof⟩

lemma *shiftr1-bl-of*:

$\text{length } bl \leq \text{len-of TYPE } ('a) \implies$
 $\text{shiftr1 } (\text{of-bl } bl :: 'a :: \text{len0 word}) = \text{of-bl } (\text{butlast } bl)$
 ⟨proof⟩

lemma *shiftr-bl-of*:

$\text{length } bl \leq \text{len-of TYPE } ('a) \implies$
 $(\text{of-bl } bl :: 'a :: \text{len0 word}) \gg n = \text{of-bl } (\text{take } (\text{length } bl - n) \ bl)$
 ⟨proof⟩

lemma *shiftr-bl*:

$(x :: 'a :: \text{len0 word}) \gg n \equiv \text{of-bl } (\text{take } (\text{len-of TYPE } ('a) - n) \ (\text{to-bl } x))$
 ⟨proof⟩

lemma *msb-shift*:

$msb (w :: 'a :: len \text{ word}) \longleftrightarrow (w \gg (len\text{-of } TYPE('a) - 1)) \neq 0$
 ⟨proof⟩

lemma *zip-replicate*:

$n \geq length \text{ ys} \implies zip (replicate \ n \ x) \text{ ys} = map (\lambda y. (x, y)) \text{ ys}$
 ⟨proof⟩

lemma *align-lem-or* [rule-format] :

$ALL \ x \ m. length \ x = n + m \dashrightarrow length \ y = n + m \dashrightarrow$
 $drop \ m \ x = replicate \ n \ False \dashrightarrow take \ m \ y = replicate \ m \ False \dashrightarrow$
 $map2 \ op \ | \ x \ y = take \ m \ x \ @ \ drop \ m \ y$
 ⟨proof⟩

lemma *align-lem-and* [rule-format] :

$ALL \ x \ m. length \ x = n + m \dashrightarrow length \ y = n + m \dashrightarrow$
 $drop \ m \ x = replicate \ n \ False \dashrightarrow take \ m \ y = replicate \ m \ False \dashrightarrow$
 $map2 \ op \ \& \ x \ y = replicate \ (n + m) \ False$
 ⟨proof⟩

lemma *aligned-bl-add-size* [OF refl]:

$size \ x - n = m \implies n \leq size \ x \implies drop \ m \ (to\text{-bl} \ x) = replicate \ n \ False \implies$
 $take \ m \ (to\text{-bl} \ y) = replicate \ m \ False \implies$
 $to\text{-bl} \ (x + y) = take \ m \ (to\text{-bl} \ x) \ @ \ drop \ m \ (to\text{-bl} \ y)$
 ⟨proof⟩

16.25.2 Mask

lemma *nth-mask* [OF refl, simp]:

$m = mask \ n \implies test\text{-bit} \ m \ i = (i < n \ \& \ i < size \ m)$
 ⟨proof⟩

lemma *mask-bl*: $mask \ n = of\text{-bl} \ (replicate \ n \ True)$

⟨proof⟩

lemma *mask-bin*: $mask \ n = word\text{-of-int} \ (bintrunc \ n \ (- \ 1))$

⟨proof⟩

lemma *and-mask-bintr*: $w \ AND \ mask \ n = word\text{-of-int} \ (bintrunc \ n \ (uint \ w))$

⟨proof⟩

lemma *and-mask-wi*: $word\text{-of-int} \ i \ AND \ mask \ n = word\text{-of-int} \ (bintrunc \ n \ i)$

⟨proof⟩

lemma *and-mask-no*: $numeral \ i \ AND \ mask \ n = word\text{-of-int} \ (bintrunc \ n \ (numeral \ i))$

⟨proof⟩

lemma *bl-and-mask'*:

$to\text{-bl} \ (w \ AND \ mask \ n :: 'a :: len \ \text{word}) =$

replicate (*len-of TYPE('a) - n*) *False* @
drop (*len-of TYPE('a) - n*) (*to-bl w*)
 ⟨*proof*⟩

lemma *and-mask-mod-2p*: $w \text{ AND mask } n = \text{word-of-int } (\text{uint } w \text{ mod } 2 \wedge n)$
 ⟨*proof*⟩

lemma *and-mask-lt-2p*: $\text{uint } (w \text{ AND mask } n) < 2 \wedge n$
 ⟨*proof*⟩

lemma *eq-mod-iff*: $0 < (n::\text{int}) \implies b = b \text{ mod } n \iff 0 \leq b \wedge b < n$
 ⟨*proof*⟩

lemma *mask-eq-iff*: $(w \text{ AND mask } n) = w \iff \text{uint } w < 2 \wedge n$
 ⟨*proof*⟩

lemma *and-mask-dvd*: $2 \wedge n \text{ dvd uint } w = (w \text{ AND mask } n = 0)$
 ⟨*proof*⟩

lemma *and-mask-dvd-nat*: $2 \wedge n \text{ dvd unat } w = (w \text{ AND mask } n = 0)$
 ⟨*proof*⟩

lemma *word-2p-lem*:
 $n < \text{size } w \implies w < 2 \wedge n = (\text{uint } (w :: 'a :: \text{len word}) < 2 \wedge n)$
 ⟨*proof*⟩

lemma *less-mask-eq*: $x < 2 \wedge n \implies x \text{ AND mask } n = (x :: 'a :: \text{len word})$
 ⟨*proof*⟩

lemmas *mask-eq-iff-w2p* = *trans* [*OF mask-eq-iff word-2p-lem* [*symmetric*]]

lemmas *and-mask-less'* = *iffD2* [*OF word-2p-lem and-mask-lt-2p, simplified word-size*]

lemma *and-mask-less-size*: $n < \text{size } x \implies x \text{ AND mask } n < 2 \wedge n$
 ⟨*proof*⟩

lemma *word-mod-2p-is-mask* [*OF refl*]:
 $c = 2 \wedge n \implies c > 0 \implies x \text{ mod } c = (x :: 'a :: \text{len word}) \text{ AND mask } n$
 ⟨*proof*⟩

lemma *mask-eqs*:
 $(a \text{ AND mask } n) + b \text{ AND mask } n = a + b \text{ AND mask } n$
 $a + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - b \text{ AND mask } n = a - b \text{ AND mask } n$
 $a - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$
 $a * (b \text{ AND mask } n) \text{ AND mask } n = a * b \text{ AND mask } n$
 $(b \text{ AND mask } n) * a \text{ AND mask } n = b * a \text{ AND mask } n$
 $(a \text{ AND mask } n) + (b \text{ AND mask } n) \text{ AND mask } n = a + b \text{ AND mask } n$
 $(a \text{ AND mask } n) - (b \text{ AND mask } n) \text{ AND mask } n = a - b \text{ AND mask } n$

$(a \text{ AND } \text{mask } n) * (b \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = a * b \text{ AND } \text{mask } n$
 $-(a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = -a \text{ AND } \text{mask } n$
 $\text{word-succ } (a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = \text{word-succ } a \text{ AND } \text{mask } n$
 $\text{word-pred } (a \text{ AND } \text{mask } n) \text{ AND } \text{mask } n = \text{word-pred } a \text{ AND } \text{mask } n$
 <proof>

lemma *mask-power-eq*:

$(x \text{ AND } \text{mask } n) \wedge k \text{ AND } \text{mask } n = x \wedge k \text{ AND } \text{mask } n$
 <proof>

16.25.3 Recast

lemmas *recast-def'* = *recast-def* [*simplified*]

lemmas *recast-def''* = *recast-def'* [*simplified word-size*]

lemmas *recast-no-def* [*simp*] = *recast-def'* [**where** *w=numeral w, unfolded word-size*]
for *w*

lemma *to-bl-recast*:

$\text{to-bl } (\text{recast } w :: 'a :: \text{len0 word}) =$
 $\text{takefill False } (\text{len-of TYPE } ('a)) (\text{to-bl } w)$
 <proof>

lemma *recast-rev-ucast* [*OF refl refl refl*]:

$cs = [rc, uc] \implies rc = \text{recast } (\text{word-reverse } w) \implies uc = \text{ucast } w \implies$
 $rc = \text{word-reverse } uc$
 <proof>

lemma *recast-ucast*: $\text{recast } w = \text{word-reverse } (\text{ucast } (\text{word-reverse } w))$
 <proof>

lemma *ucast-recast*: $\text{ucast } w = \text{word-reverse } (\text{recast } (\text{word-reverse } w))$
 <proof>

lemma *ucast-rev-recast*: $\text{ucast } (\text{word-reverse } w) = \text{word-reverse } (\text{recast } w)$
 <proof>

lemmas *wsst-TYs* = *source-size target-size word-size*

lemma *recast-down-uu* [*OF refl*]:

$rc = \text{recast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc (w :: 'a :: \text{len word}) = \text{ucast } (w \gg n)$
 <proof>

lemma *recast-down-us* [*OF refl*]:

$rc = \text{recast} \implies \text{source-size } rc = \text{target-size } rc + n \implies$
 $rc (w :: 'a :: \text{len word}) = \text{ucast } (w \gg \gg n)$
 <proof>

lemma *recast-down-su* [*OF refl*]:

$rc = revcast \implies source\text{-}size\ rc = target\text{-}size\ rc + n \implies$
 $rc\ (w :: 'a :: len\ word) = scast\ (w >> n)$
 ⟨proof⟩

lemma *revcast-down-ss* [OF refl]:

$rc = revcast \implies source\text{-}size\ rc = target\text{-}size\ rc + n \implies$
 $rc\ (w :: 'a :: len\ word) = scast\ (w >>> n)$
 ⟨proof⟩

lemma *cast-down-rev*:

$uc = ucast \implies source\text{-}size\ uc = target\text{-}size\ uc + n \implies$
 $uc\ w = revcast\ ((w :: 'a :: len\ word) << n)$
 ⟨proof⟩

lemma *revcast-up* [OF refl]:

$rc = revcast \implies source\text{-}size\ rc + n = target\text{-}size\ rc \implies$
 $rc\ w = (ucast\ w :: 'a :: len\ word) << n$
 ⟨proof⟩

lemmas *rc1 = revcast-up* [THEN

revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

lemmas *rc2 = revcast-down-uu* [THEN

revcast-rev-ucast [symmetric, THEN trans, THEN word-rev-gal, symmetric]]

lemmas *ucast-up =*

rc1 [simplified rev-shiftr [symmetric] revcast-ucast [symmetric]]

lemmas *ucast-down =*

rc2 [simplified rev-shiftr revcast-ucast [symmetric]]

16.25.4 Slices

lemma *slice1-no-bin* [simp]:

$slice1\ n\ (numeral\ w :: 'b\ word) = of\text{-}bl\ (takefill\ False\ n\ (bin\text{-}to\text{-}bl\ (len\text{-}of\ TYPE('b$
 $:: len0)))\ (numeral\ w)))$
 ⟨proof⟩

lemma *slice-no-bin* [simp]:

$slice\ n\ (numeral\ w :: 'b\ word) = of\text{-}bl\ (takefill\ False\ (len\text{-}of\ TYPE('b :: len0) -$
 $n)$
 $(bin\text{-}to\text{-}bl\ (len\text{-}of\ TYPE('b :: len0)))\ (numeral\ w)))$
 ⟨proof⟩

lemma *slice1-0* [simp] : $slice1\ n\ 0 = 0$

⟨proof⟩

lemma *slice-0* [simp] : $slice\ n\ 0 = 0$

⟨proof⟩

lemma *slice-take'*: $\text{slice } n \ w = \text{of-bl } (\text{take } (\text{size } w - n) \ (\text{to-bl } w))$
 ⟨proof⟩

lemmas *slice-take* = *slice-take'* [unfolded word-size]

— shiftr to a word of the same size is just slice, slice is just shiftr then ucast

lemmas *shiftr-slice* = *trans* [OF *shiftr-bl* [THEN *meta-eq-to-obj-eq*] *slice-take* [symmetric]]

lemma *slice-shiftr*: $\text{slice } n \ w = \text{ucast } (w \gg n)$
 ⟨proof⟩

lemma *nth-slice*:
 $(\text{slice } n \ w :: 'a :: \text{len0 word}) !! m =$
 $(w !! (m + n) \ \& \ m < \text{len-of TYPE } ('a))$
 ⟨proof⟩

lemma *slice1-down-alt'*:
 $sl = \text{slice1 } n \ w \implies fs = \text{size } sl \implies fs + k = n \implies$
 $\text{to-bl } sl = \text{takefill False } fs \ (\text{drop } k \ (\text{to-bl } w))$
 ⟨proof⟩

lemma *slice1-up-alt'*:
 $sl = \text{slice1 } n \ w \implies fs = \text{size } sl \implies fs = n + k \implies$
 $\text{to-bl } sl = \text{takefill False } fs \ (\text{replicate } k \ \text{False } @ \ (\text{to-bl } w))$
 ⟨proof⟩

lemmas *sd1* = *slice1-down-alt'* [OF *refl refl*, unfolded word-size]

lemmas *su1* = *slice1-up-alt'* [OF *refl refl*, unfolded word-size]

lemmas *slice1-down-alt* = *le-add-diff-inverse* [THEN *sd1*]

lemmas *slice1-up-alt* =
le-add-diff-inverse [symmetric, THEN *su1*]
le-add-diff-inverse2 [symmetric, THEN *su1*]

lemma *ucast-slice1*: $\text{ucast } w = \text{slice1 } (\text{size } w) \ w$
 ⟨proof⟩

lemma *ucast-slice*: $\text{ucast } w = \text{slice } 0 \ w$
 ⟨proof⟩

lemma *slice-id*: $\text{slice } 0 \ t = t$
 ⟨proof⟩

lemma *revcast-slice1* [OF *refl*]:
 $rc = \text{revcast } w \implies \text{slice1 } (\text{size } rc) \ w = rc$
 ⟨proof⟩

lemma *slice1-tf-tf'*:
 $\text{to-bl } (\text{slice1 } n \ w :: 'a :: \text{len0 word}) =$
 $\text{rev } (\text{takefill False } (\text{len-of TYPE } ('a)) \ (\text{rev } (\text{takefill False } n \ (\text{to-bl } w))))$

<proof>

lemmas *slice1-tf-tf* = *slice1-tf-tf'* [*THEN word-bl.Rep-inverse'*, *symmetric*]

lemma *rev-slice1*:

$n + k = \text{len-of TYPE}('a) + \text{len-of TYPE}('b) \implies$
 $\text{slice1 } n (\text{word-reverse } w :: 'b :: \text{len0 word}) =$
 $\text{word-reverse } (\text{slice1 } k w :: 'a :: \text{len0 word})$
<proof>

lemma *rev-slice*:

$n + k + \text{len-of TYPE}('a::\text{len0}) = \text{len-of TYPE}('b::\text{len0}) \implies$
 $\text{slice } n (\text{word-reverse } (w::'b \text{ word})) = \text{word-reverse } (\text{slice } k w::'a \text{ word})$
<proof>

lemmas *sym-notr* =

not-iff [*THEN iffD2*, *THEN not-sym*, *THEN not-iff* [*THEN iffD1*]]

— problem posed by TPHOLs referee: criterion for overflow of addition of signed integers

lemma *soft-test*:

$(\text{sint } (x :: 'a :: \text{len word}) + \text{sint } y = \text{sint } (x + y)) =$
 $((((x+y) \text{ XOR } x) \text{ AND } ((x+y) \text{ XOR } y)) \gg (\text{size } x - 1) = 0)$
<proof>

16.26 Split and cat

lemmas *word-split-bin'* = *word-split-def*

lemmas *word-cat-bin'* = *word-cat-def*

lemma *word-rsplit-no*:

$(\text{word-rsplit } (\text{numeral bin} :: 'b :: \text{len0 word}) :: 'a \text{ word list}) =$
 $\text{map word-of-int } (\text{bin-rsplit } (\text{len-of TYPE}('a :: \text{len}))$
 $(\text{len-of TYPE}('b), \text{bintrunc } (\text{len-of TYPE}('b)) (\text{numeral bin})))$
<proof>

lemmas *word-rsplit-no-cl* [*simp*] = *word-rsplit-no*

[*unfolded bin-rsplitl-def bin-rsplit-l* [*symmetric*]]

lemma *test-bit-cat*:

$wc = \text{word-cat } a \ b \implies wc !! n = (n < \text{size } wc \ \&$
 $(\text{if } n < \text{size } b \ \text{then } b !! n \ \text{else } a !! (n - \text{size } b)))$
<proof>

lemma *word-cat-bl*: $\text{word-cat } a \ b = \text{of-bl } (\text{to-bl } a \ @ \ \text{to-bl } b)$

<proof>

lemma *of-bl-append*:

$(\text{of-bl } (xs \text{ @ } ys) :: 'a :: \text{len word}) = \text{of-bl } xs * 2^{(\text{length } ys)} + \text{of-bl } ys$
 ⟨proof⟩

lemma *of-bl-False* [simp]:
 $\text{of-bl } (\text{False}\#xs) = \text{of-bl } xs$
 ⟨proof⟩

lemma *of-bl-True* [simp]:
 $(\text{of-bl } (\text{True}\#xs)::'a::\text{len word}) = 2^{\text{length } xs} + \text{of-bl } xs$
 ⟨proof⟩

lemma *of-bl-Cons*:
 $\text{of-bl } (x\#xs) = \text{of-bool } x * 2^{\text{length } xs} + \text{of-bl } xs$
 ⟨proof⟩

lemma *split-wint-lem*: $\text{bin-split } n \ (\text{uint } (w :: 'a :: \text{len0 word})) = (a, b) \implies$
 $a = \text{bintrunc } (\text{len-of TYPE('a)} - n) \ a \ \& \ b = \text{bintrunc } (\text{len-of TYPE('a)}) \ b$
 ⟨proof⟩

lemma *word-split-bl'*:
 $\text{std} = \text{size } c - \text{size } b \implies (\text{word-split } c = (a, b)) \implies$
 $(a = \text{of-bl } (\text{take } \text{std } (\text{to-bl } c)) \ \& \ b = \text{of-bl } (\text{drop } \text{std } (\text{to-bl } c)))$
 ⟨proof⟩

lemma *word-split-bl*: $\text{std} = \text{size } c - \text{size } b \implies$
 $(a = \text{of-bl } (\text{take } \text{std } (\text{to-bl } c)) \ \& \ b = \text{of-bl } (\text{drop } \text{std } (\text{to-bl } c))) \longleftrightarrow$
 $\text{word-split } c = (a, b)$
 ⟨proof⟩

lemma *word-split-bl-eq*:
 $(\text{word-split } (c::'a::\text{len word}) :: ('c :: \text{len0 word} * 'd :: \text{len0 word})) =$
 $(\text{of-bl } (\text{take } (\text{len-of TYPE('a)::len}) - \text{len-of TYPE('d)::len0})) (\text{to-bl } c),$
 $\text{of-bl } (\text{drop } (\text{len-of TYPE('a)} - \text{len-of TYPE('d)}) (\text{to-bl } c))$
 ⟨proof⟩

lemma *test-bit-split'*:
 $\text{word-split } c = (a, b) \dashrightarrow (\text{ALL } n \ m. \ b \ !! \ n = (n < \text{size } b \ \& \ c \ !! \ n) \ \&$
 $a \ !! \ m = (m < \text{size } a \ \& \ c \ !! \ (m + \text{size } b)))$
 ⟨proof⟩

lemma *test-bit-split*:
 $\text{word-split } c = (a, b) \implies$
 $(\forall n::\text{nat}. \ b \ !! \ n \longleftrightarrow n < \text{size } b \ \wedge \ c \ !! \ n) \ \wedge \ (\forall m::\text{nat}. \ a \ !! \ m \longleftrightarrow m < \text{size } a$
 $\wedge \ c \ !! \ (m + \text{size } b))$
 ⟨proof⟩

lemma *test-bit-split-eq*: $\text{word-split } c = (a, b) \longleftrightarrow$
 $((\text{ALL } n::\text{nat}. \ b \ !! \ n = (n < \text{size } b \ \& \ c \ !! \ n)) \ \&$
 $(\text{ALL } m::\text{nat}. \ a \ !! \ m = (m < \text{size } a \ \& \ c \ !! \ (m + \text{size } b))))$
 ⟨proof⟩

lemma *word-cat-id*: $\text{word-cat } a \ b = b$

<proof>

lemma *word-cat-hom*:

$\text{len-of TYPE}('a::\text{len0}) \leq \text{len-of TYPE}('b::\text{len0}) + \text{len-of TYPE}('c::\text{len0})$

\implies

$(\text{word-cat } (\text{word-of-int } w :: 'b \ \text{word}) \ (b :: 'c \ \text{word}) :: 'a \ \text{word}) =$

$\text{word-of-int } (\text{bin-cat } w \ (\text{size } b) \ (\text{uint } b))$

<proof>

lemma *word-cat-split-alt*:

$\text{size } w \leq \text{size } u + \text{size } v \implies \text{word-split } w = (u, v) \implies \text{word-cat } u \ v = w$

<proof>

lemmas *word-cat-split-size = sym [THEN [2] word-cat-split-alt [symmetric]]*

16.26.1 Split and slice

lemma *split-slices*:

$\text{word-split } w = (u, v) \implies u = \text{slice } (\text{size } v) \ w \ \& \ v = \text{slice } 0 \ w$

<proof>

lemma *slice-cat1 [OF refl]*:

$wc = \text{word-cat } a \ b \implies \text{size } wc \geq \text{size } a + \text{size } b \implies \text{slice } (\text{size } b) \ wc = a$

<proof>

lemmas *slice-cat2 = trans [OF slice-id word-cat-id]*

lemma *cat-slices*:

$a = \text{slice } n \ c \implies b = \text{slice } 0 \ c \implies n = \text{size } b \implies$

$\text{size } a + \text{size } b \geq \text{size } c \implies \text{word-cat } a \ b = c$

<proof>

lemma *word-split-cat-alt*:

$w = \text{word-cat } u \ v \implies \text{size } u + \text{size } v \leq \text{size } w \implies \text{word-split } w = (u, v)$

<proof>

lemmas *word-cat-bl-no-bin [simp] =*

word-cat-bl [where a=numeral a and b=numeral b,

unfolded to-bl-numeral]

for $a \ b$

lemmas *word-split-bl-no-bin [simp] =*

word-split-bl-eq [where c=numeral c, unfolded to-bl-numeral] for c

this odd result arises from the fact that the statement of the result implies that the decoded words are of the same type, and therefore of the same length, as the original word

lemma *word-rsplit-same*: $\text{word-rsplit } w = [w]$

<proof>

lemma *word-rsplit-empty-iff-size*:
 $(\text{word-rsplit } w = []) = (\text{size } w = 0)$
<proof>

lemma *test-bit-rsplit*:
 $sw = \text{word-rsplit } w \implies m < \text{size } (\text{hd } sw :: 'a :: \text{len } \text{word}) \implies$
 $k < \text{length } sw \implies (\text{rev } sw ! k) !! m = (w !! (k * \text{size } (\text{hd } sw) + m))$
<proof>

lemma *word-rcat-bl*: $\text{word-rcat } wl = \text{of-bl } (\text{concat } (\text{map } \text{to-bl } wl))$
<proof>

lemma *size-rcat-lem'*:
 $\text{size } (\text{concat } (\text{map } \text{to-bl } wl)) = \text{length } wl * \text{size } (\text{hd } wl)$
<proof>

lemmas *size-rcat-lem = size-rcat-lem'* [*unfolded word-size*]

lemmas *td-gal-lt-len = len-gt-0* [*THEN td-gal-lt*]

lemma *nth-rcat-lem*:
 $n < \text{length } (wl :: 'a \text{ word list}) * \text{len-of } \text{TYPE}('a :: \text{len}) \implies$
 $\text{rev } (\text{concat } (\text{map } \text{to-bl } wl)) ! n =$
 $\text{rev } (\text{to-bl } (\text{rev } wl ! (n \text{ div } \text{len-of } \text{TYPE}('a)))) ! (n \text{ mod } \text{len-of } \text{TYPE}('a))$
<proof>

lemma *test-bit-rcat*:
 $sw = \text{size } (\text{hd } wl :: 'a :: \text{len } \text{word}) \implies rc = \text{word-rcat } wl \implies rc !! n =$
 $(n < \text{size } rc \ \& \ n \text{ div } sw < \text{size } wl \ \& \ (\text{rev } wl) ! (n \text{ div } sw) !! (n \text{ mod } sw))$
<proof>

lemma *foldl-eq-foldr*:
 $\text{foldl } op + x \ xs = \text{foldr } op + (x \ \# \ xs) \ (0 :: 'a :: \text{comm-monoid-add})$
<proof>

lemmas *test-bit-cong = arg-cong* [**where** $f = \text{test-bit}$, *THEN fun-cong*]

lemmas *test-bit-rsplit-alt =*
 $\text{trans } [\text{OF } \text{nth-rev-alt } [\text{THEN } \text{test-bit-cong}]]$
 $\text{test-bit-rsplit } [\text{OF } \text{refl } \text{asm-rl } \text{diff-Suc-less}]$

— lazy way of expressing that u and v , and su and sv , have same types

lemma *word-rsplit-len-indep* [*OF refl refl refl refl*]:
 $[u, v] = p \implies [su, sv] = q \implies \text{word-rsplit } u = su \implies$
 $\text{word-rsplit } v = sv \implies \text{length } su = \text{length } sv$
<proof>

lemma *length-word-rsplit-size*:

$n = \text{len-of TYPE } ('a :: \text{len}) \implies$
 $(\text{length } (\text{word-rsplit } w :: 'a \text{ word list}) \leq m) = (\text{size } w \leq m * n)$
 ⟨proof⟩

lemmas *length-word-rsplit-lt-size* =

length-word-rsplit-size [unfolded Not-eq-iff linorder-not-less [symmetric]]

lemma *length-word-rsplit-exp-size*:

$n = \text{len-of TYPE } ('a :: \text{len}) \implies$
 $\text{length } (\text{word-rsplit } w :: 'a \text{ word list}) = (\text{size } w + n - 1) \text{ div } n$
 ⟨proof⟩

lemma *length-word-rsplit-even-size*:

$n = \text{len-of TYPE } ('a :: \text{len}) \implies \text{size } w = m * n \implies$
 $\text{length } (\text{word-rsplit } w :: 'a \text{ word list}) = m$
 ⟨proof⟩

lemmas *length-word-rsplit-exp-size'* = refl [THEN *length-word-rsplit-exp-size*]

lemmas *tdle* = iffD2 [OF *split-div-lemma* refl, THEN *conjunct1*]

lemmas *dtle* = xtr4 [OF *tdle* mult.commute]

lemma *word-rcat-rsplit*: $\text{word-rcat } (\text{word-rsplit } w) = w$
 ⟨proof⟩

lemma *size-word-rsplit-rcat-size*:

$\llbracket \text{word-rcat } (ws :: 'a :: \text{len word list}) = (\text{frcw} :: 'b :: \text{len0 word});$
 $\text{size frcw} = \text{length } ws * \text{len-of TYPE } ('a) \rrbracket$
 $\implies \text{length } (\text{word-rsplit } \text{frcw} :: 'a \text{ word list}) = \text{length } ws$
 ⟨proof⟩

lemma *msreus*:

fixes $n :: \text{nat}$
shows $0 < n \implies (k * n + m) \text{ div } n = m \text{ div } n + k$
and $(k * n + m) \text{ mod } n = m \text{ mod } n$
 ⟨proof⟩

lemma *word-rsplit-rcat-size* [OF refl]:

$\text{word-rcat } (ws :: 'a :: \text{len word list}) = \text{frcw} \implies$
 $\text{size frcw} = \text{length } ws * \text{len-of TYPE } ('a) \implies \text{word-rsplit } \text{frcw} = ws$
 ⟨proof⟩

16.27 Rotation

lemmas *rotater-0'* [simp] = *rotater-def* [where $n = 0$, simplified]

lemmas *word-rot-defs* = *word-roti-def* *word-rotr-def* *word-rotl-def*

lemma *rotate-eq-mod*:

$m \bmod \text{length } xs = n \bmod \text{length } xs \implies \text{rotate } m \text{ } xs = \text{rotate } n \text{ } xs$
 ⟨proof⟩

lemmas *rotate-eqs* =

trans [OF rotate0 [THEN fun-cong] id-apply]
rotate-rotate [symmetric]
rotate-id
rotate-conv-mod
rotate-eq-mod

16.27.1 Rotation of list to right

lemma *rotate1-rl'*: $\text{rotater1 } (l @ [a]) = a \# l$
 ⟨proof⟩

lemma *rotate1-rl* [simp]: $\text{rotater1 } (\text{rotate1 } l) = l$
 ⟨proof⟩

lemma *rotate1-lr* [simp]: $\text{rotate1 } (\text{rotater1 } l) = l$
 ⟨proof⟩

lemma *rotater1-rev'*: $\text{rotater1 } (\text{rev } xs) = \text{rev } (\text{rotate1 } xs)$
 ⟨proof⟩

lemma *rotater-rev'*: $\text{rotater } n (\text{rev } xs) = \text{rev } (\text{rotate } n \text{ } xs)$
 ⟨proof⟩

lemma *rotater-rev*: $\text{rotater } n \text{ } ys = \text{rev } (\text{rotate } n (\text{rev } ys))$
 ⟨proof⟩

lemma *rotater-drop-take*:

$\text{rotater } n \text{ } xs =$
 $\text{drop } (\text{length } xs - n \bmod \text{length } xs) \text{ } xs @$
 $\text{take } (\text{length } xs - n \bmod \text{length } xs) \text{ } xs$
 ⟨proof⟩

lemma *rotater-Suc* [simp]:

$\text{rotater } (\text{Suc } n) \text{ } xs = \text{rotater1 } (\text{rotater } n \text{ } xs)$
 ⟨proof⟩

lemma *rotate-inv-plus* [rule-format]:

ALL k . $k = m + n \implies \text{rotater } k (\text{rotate } n \text{ } xs) = \text{rotater } m \text{ } xs \ \&$
 $\text{rotate } k (\text{rotater } n \text{ } xs) = \text{rotate } m \text{ } xs \ \&$
 $\text{rotater } n (\text{rotate } k \text{ } xs) = \text{rotate } m \text{ } xs \ \&$
 $\text{rotate } n (\text{rotater } k \text{ } xs) = \text{rotater } m \text{ } xs$
 ⟨proof⟩

lemmas *rotate-inv-rel* = *le-add-diff-inverse2* [*symmetric*, *THEN rotate-inv-plus*]

lemmas *rotate-inv-eq* = *order-refl* [*THEN rotate-inv-rel*, *simplified*]

lemmas *rotate-lr* [*simp*] = *rotate-inv-eq* [*THEN conjunct1*]

lemmas *rotate-rl* [*simp*] = *rotate-inv-eq* [*THEN conjunct2*, *THEN conjunct1*]

lemma *rotate-gal*: $(\text{rotater } n \text{ } xs = ys) = (\text{rotate } n \text{ } ys = xs)$
 ⟨*proof*⟩

lemma *rotate-gal'*: $(ys = \text{rotater } n \text{ } xs) = (xs = \text{rotate } n \text{ } ys)$
 ⟨*proof*⟩

lemma *length-rotater* [*simp*]:
 $\text{length } (\text{rotater } n \text{ } xs) = \text{length } xs$
 ⟨*proof*⟩

lemma *restrict-to-left*:
assumes $x = y$
shows $(x = z) = (y = z)$
 ⟨*proof*⟩

lemmas *rrs0* = *rotate-eqs* [*THEN restrict-to-left*,
simplified rotate-gal [*symmetric*] *rotate-gal'* [*symmetric*]]
lemmas *rrs1* = *rrs0* [*THEN refl* [*THEN rev-iffD1*]]
lemmas *rotater-eqs* = *rrs1* [*simplified length-rotater*]
lemmas *rotater-0* = *rotater-eqs* (1)
lemmas *rotater-add* = *rotater-eqs* (2)

16.27.2 map, map2, commuting with rotate(r)

lemma *butlast-map*:
 $xs \sim = [] \implies \text{butlast } (\text{map } f \text{ } xs) = \text{map } f \text{ } (\text{butlast } xs)$
 ⟨*proof*⟩

lemma *rotater1-map*: $\text{rotater1 } (\text{map } f \text{ } xs) = \text{map } f \text{ } (\text{rotater1 } xs)$
 ⟨*proof*⟩

lemma *rotater-map*:
 $\text{rotater } n \text{ } (\text{map } f \text{ } xs) = \text{map } f \text{ } (\text{rotater } n \text{ } xs)$
 ⟨*proof*⟩

lemma *but-last-zip* [*rule-format*] :
 $ALL \text{ } ys. \text{length } xs = \text{length } ys \implies xs \sim = [] \implies$
 $\text{last } (\text{zip } xs \text{ } ys) = (\text{last } xs, \text{last } ys) \ \&$
 $\text{butlast } (\text{zip } xs \text{ } ys) = \text{zip } (\text{butlast } xs) \text{ } (\text{butlast } ys)$
 ⟨*proof*⟩

lemma *but-last-map2* [*rule-format*] :

$ALL\ ys.\ length\ xs = length\ ys \dashrightarrow xs \sim = [] \dashrightarrow$
 $last\ (map2\ f\ xs\ ys) = f\ (last\ xs)\ (last\ ys)\ \&$
 $butlast\ (map2\ f\ xs\ ys) = map2\ f\ (butlast\ xs)\ (butlast\ ys)$
 <proof>

lemma *rotater1-zip*:

$length\ xs = length\ ys \implies$
 $rotater1\ (zip\ xs\ ys) = zip\ (rotater1\ xs)\ (rotater1\ ys)$
 <proof>

lemma *rotater1-map2*:

$length\ xs = length\ ys \implies$
 $rotater1\ (map2\ f\ xs\ ys) = map2\ f\ (rotater1\ xs)\ (rotater1\ ys)$
 <proof>

lemmas *lrth* =

$box-equals\ [OF\ asm-rl\ length-rotater\ [symmetric]$
 $length-rotater\ [symmetric],$
 $THEN\ rotater1-map2]$

lemma *rotater-map2*:

$length\ xs = length\ ys \implies$
 $rotater\ n\ (map2\ f\ xs\ ys) = map2\ f\ (rotater\ n\ xs)\ (rotater\ n\ ys)$
 <proof>

lemma *rotate1-map2*:

$length\ xs = length\ ys \implies$
 $rotate1\ (map2\ f\ xs\ ys) = map2\ f\ (rotate1\ xs)\ (rotate1\ ys)$
 <proof>

lemmas *lth* = $box-equals\ [OF\ asm-rl\ length-rotate\ [symmetric]$

$length-rotate\ [symmetric],\ THEN\ rotate1-map2]$

lemma *rotate-map2*:

$length\ xs = length\ ys \implies$
 $rotate\ n\ (map2\ f\ xs\ ys) = map2\ f\ (rotate\ n\ xs)\ (rotate\ n\ ys)$
 <proof>

lemma *to-bl-rotl*:

$to-bl\ (word-rotl\ n\ w) = rotate\ n\ (to-bl\ w)$
 <proof>

lemmas *blrs0* = $rotate-egs\ [THEN\ to-bl-rotl\ [THEN\ trans]]$

lemmas *word-rotl-egs* =

$blrs0\ [simplified\ word-bl-Rep'\ word-bl.Rep-inject\ to-bl-rotl\ [symmetric]]$

lemma *to-bl-rotr*:

$to-bl\ (word-rotr\ n\ w) = rotater\ n\ (to-bl\ w)$

<proof>

lemmas *brrs0 = rotater-eqs* [*THEN to-bl-rotr* [*THEN trans*]]

lemmas *word-rotr-eqs =*

brrs0 [*simplified word-bl-Rep'* *word-bl.Rep-inject to-bl-rotr* [*symmetric*]]

declare *word-rotr-eqs* (1) [*simp*]

declare *word-rotl-eqs* (1) [*simp*]

lemma

word-rot-rl [*simp*]:

word-rotl k (word-rotr k v) = v **and**

word-rot-lr [*simp*]:

word-rotr k (word-rotl k v) = v

<proof>

lemma

word-rot-gal:

(word-rotr n v = w) = (word-rotl n w = v) **and**

word-rot-gal':

(w = word-rotr n v) = (v = word-rotl n w)

<proof>

lemma *word-rotr-rev*:

word-rotr n w = word-reverse (word-rotl n (word-reverse w))

<proof>

lemma *word-roti-0* [*simp*]: *word-roti 0 w = w*

<proof>

lemmas *abl-cong = arg-cong* [**where** *f = of-bl*]

lemma *word-roti-add*:

word-roti (m + n) w = word-roti m (word-roti n w)

<proof>

lemma *word-roti-conv-mod'*: *word-roti n w = word-roti (n mod int (size w)) w*

<proof>

lemmas *word-roti-conv-mod = word-roti-conv-mod'* [*unfolded word-size*]

16.27.3 ”Word rotation commutes with bit-wise operations

locale *word-rotate*

begin

lemmas *word-rot-defs' = to-bl-rotl to-bl-rotr*

lemmas *blwl-syms* [*symmetric*] = *bl-word-not bl-word-and bl-word-or bl-word-xor*

lemmas *lbl-lbl* = *trans* [*OF word-bl-Rep' word-bl-Rep' [symmetric]*]

lemmas *ths-map2* [*OF lbl-lbl*] = *rotate-map2 rotater-map2*

lemmas *ths-map* [**where** *xs = to-bl v*] = *rotate-map rotater-map* **for** *v*

lemmas *th1s* [*simplified word-rot-defs' [symmetric]*] = *ths-map2 ths-map*

lemma *word-rot-logs*:

word-rotl n (NOT v) = NOT word-rotl n v

word-rotr n (NOT v) = NOT word-rotr n v

word-rotl n (x AND y) = word-rotl n x AND word-rotl n y

word-rotr n (x AND y) = word-rotr n x AND word-rotr n y

word-rotl n (x OR y) = word-rotl n x OR word-rotl n y

word-rotr n (x OR y) = word-rotr n x OR word-rotr n y

word-rotl n (x XOR y) = word-rotl n x XOR word-rotl n y

word-rotr n (x XOR y) = word-rotr n x XOR word-rotr n y

<proof>

end

lemmas *word-rot-logs* = *word-rotate.word-rot-logs*

lemmas *bl-word-rotl-dt* = *trans* [*OF to-bl-rotl rotate-drop-take, simplified word-bl-Rep'*]

lemmas *bl-word-rotr-dt* = *trans* [*OF to-bl-rotr rotater-drop-take, simplified word-bl-Rep'*]

lemma *bl-word-roti-dt'*:

n = nat ((- i) mod int (size (w :: 'a :: len word))) \implies

to-bl (word-roti i w) = drop n (to-bl w) @ take n (to-bl w)

<proof>

lemmas *bl-word-roti-dt* = *bl-word-roti-dt'* [*unfolded word-size*]

lemmas *word-rotl-dt* = *bl-word-rotl-dt* [*THEN word-bl.Rep-inverse' [symmetric]*]

lemmas *word-rotr-dt* = *bl-word-rotr-dt* [*THEN word-bl.Rep-inverse' [symmetric]*]

lemmas *word-roti-dt* = *bl-word-roti-dt* [*THEN word-bl.Rep-inverse' [symmetric]*]

lemma *word-rotx-0* [*simp*] : *word-rotr i 0 = 0 & word-rotl i 0 = 0*
<proof>

lemma *word-roti-0'* [*simp*] : *word-roti n 0 = 0*
<proof>

lemmas *word-rotr-dt-no-bin'* [*simp*] =
word-rotr-dt [**where** *w=numeral w, unfolded to-bl-numeral*] **for** *w*

lemmas *word-rotl-dt-no-bin'* [*simp*] =
word-rotl-dt [**where** *w=numeral w, unfolded to-bl-numeral*] **for** *w*

declare *word-roti-def* [*simp*]

16.28 Maximum machine word

lemma *word-int-cases*:

obtains *n* **where** $(x :: 'a::len0 \text{ word}) = \text{word-of-int } n$ **and** $0 \leq n$ **and** $n < 2^{\text{len-of TYPE('a)}}$
 ⟨*proof*⟩

lemma *word-nat-cases* [*cases type: word*]:

obtains *n* **where** $(x :: 'a::len \text{ word}) = \text{of-nat } n$ **and** $n < 2^{\text{len-of TYPE('a)}}$
 ⟨*proof*⟩

lemma *max-word-eq*: $(\text{max-word} :: 'a::len \text{ word}) = 2^{\text{len-of TYPE('a)}} - 1$
 ⟨*proof*⟩

lemma *max-word-max* [*simp,intro!*]: $n \leq \text{max-word}$
 ⟨*proof*⟩

lemma *word-of-int-2p-len*: $\text{word-of-int } (2^{\text{len-of TYPE('a)}}) = (0 :: 'a::len0 \text{ word})$
 ⟨*proof*⟩

lemma *word-pow-0*:

$(2 :: 'a::len \text{ word})^{\text{len-of TYPE('a)}} = 0$
 ⟨*proof*⟩

lemma *max-word-wrap*: $x + 1 = 0 \implies x = \text{max-word}$
 ⟨*proof*⟩

lemma *max-word-minus*:

$\text{max-word} = (-1 :: 'a::len \text{ word})$
 ⟨*proof*⟩

lemma *max-word-bl* [*simp*]:

to-bl $(\text{max-word} :: 'a::len \text{ word}) = \text{replicate } (\text{len-of TYPE('a)}) \text{ True}$
 ⟨*proof*⟩

lemma *max-test-bit* [*simp*]:

$(\text{max-word} :: 'a::len \text{ word}) !! n = (n < \text{len-of TYPE('a)})$
 ⟨*proof*⟩

lemma *word-and-max* [*simp*]:

$x \text{ AND } \text{max-word} = x$

<proof>

lemma *word-or-max* [*simp*]:
 $x \text{ OR } \text{max-word} = \text{max-word}$
<proof>

lemma *word-ao-dist2*:
 $x \text{ AND } (y \text{ OR } z) = x \text{ AND } y \text{ OR } x \text{ AND } (z :: 'a :: \text{len } 0 \text{ word})$
<proof>

lemma *word-oa-dist2*:
 $x \text{ OR } y \text{ AND } z = (x \text{ OR } y) \text{ AND } (x \text{ OR } (z :: 'a :: \text{len } 0 \text{ word}))$
<proof>

lemma *word-and-not* [*simp*]:
 $x \text{ AND } \text{NOT } x = (0 :: 'a :: \text{len } 0 \text{ word})$
<proof>

lemma *word-or-not* [*simp*]:
 $x \text{ OR } \text{NOT } x = \text{max-word}$
<proof>

lemma *word-boolean*:
 $\text{boolean } (op \text{ AND}) (op \text{ OR}) \text{ bitNOT } 0 \text{ max-word}$
<proof>

interpretation *word-bool-alg*:
 $\text{boolean } op \text{ AND } op \text{ OR } \text{bitNOT } 0 \text{ max-word}$
<proof>

lemma *word-xor-and-or*:
 $x \text{ XOR } y = x \text{ AND } \text{NOT } y \text{ OR } \text{NOT } x \text{ AND } (y :: 'a :: \text{len } 0 \text{ word})$
<proof>

interpretation *word-bool-alg*:
 $\text{boolean-xor } op \text{ AND } op \text{ OR } \text{bitNOT } 0 \text{ max-word } op \text{ XOR}$
<proof>

lemma *shiftr-x-0* [*iff*]:
 $(x :: 'a :: \text{len } 0 \text{ word}) \gg 0 = x$
<proof>

lemma *shiftr-x-0* [*simp*]:
 $(x :: 'a :: \text{len } 0 \text{ word}) \gg 0 = x$
<proof>

lemma *shiftr-1* [*simp*]:
 $(1 :: 'a :: \text{len } 0 \text{ word}) \gg n = 2^n$
<proof>

lemma *uint-lt-0* [*simp*]:

$uint\ x < 0 = False$
 ⟨*proof*⟩

lemma *shiftr1-1* [*simp*]:

$shiftr1\ (1::'a::len\ word) = 0$
 ⟨*proof*⟩

lemma *shiftr-1* [*simp*]:

$(1::'a::len\ word) >> n = (if\ n = 0\ then\ 1\ else\ 0)$
 ⟨*proof*⟩

lemma *word-less-1* [*simp*]:

$((x::'a::len\ word) < 1) = (x = 0)$
 ⟨*proof*⟩

lemma *to-bl-mask*:

$to-bl\ (mask\ n :: 'a::len\ word) =$
 $replicate\ (len-of\ TYPE('a) - n)\ False\ @$
 $replicate\ (min\ (len-of\ TYPE('a))\ n)\ True$
 ⟨*proof*⟩

lemma *map-replicate-True*:

$n = length\ xs \implies$
 $map\ (\lambda(x,y).\ x\ \&\ y)\ (zip\ xs\ (replicate\ n\ True)) = xs$
 ⟨*proof*⟩

lemma *map-replicate-False*:

$n = length\ xs \implies map\ (\lambda(x,y).\ x\ \&\ y)$
 $(zip\ xs\ (replicate\ n\ False)) = replicate\ n\ False$
 ⟨*proof*⟩

lemma *bl-and-mask*:

fixes $w :: 'a::len\ word$
fixes n
defines $n' \equiv len-of\ TYPE('a) - n$
shows $to-bl\ (w\ AND\ mask\ n) = replicate\ n'\ False\ @\ drop\ n'\ (to-bl\ w)$
 ⟨*proof*⟩

lemma *drop-rev-takefill*:

$length\ xs \leq n \implies$
 $drop\ (n - length\ xs)\ (rev\ (takefill\ False\ n\ (rev\ xs))) = xs$
 ⟨*proof*⟩

lemma *map-nth-0* [*simp*]:

$map\ (op\ !!\ (0::'a::len0\ word))\ xs = replicate\ (length\ xs)\ False$
 ⟨*proof*⟩

lemma *uint-plus-if-size*:

$uint\ (x + y) =$
(if $uint\ x + uint\ y < 2^{size\ x}$ *then*
 $uint\ x + uint\ y$
else
 $uint\ x + uint\ y - 2^{size\ x}$)
 ⟨*proof*⟩

lemma *unat-plus-if-size*:

$unat\ (x + (y::'a::len\ word)) =$
(if $unat\ x + unat\ y < 2^{size\ x}$ *then*
 $unat\ x + unat\ y$
else
 $unat\ x + unat\ y - 2^{size\ x}$)
 ⟨*proof*⟩

lemma *word-neq-0-conv*:

fixes $w :: 'a :: len\ word$
shows $(w \neq 0) = (0 < w)$
 ⟨*proof*⟩

lemma *max-lt*:

$unat\ (max\ a\ b\ div\ c) = unat\ (max\ a\ b)\ div\ unat\ (c::'a::len\ word)$
 ⟨*proof*⟩

lemma *uint-sub-if-size*:

$uint\ (x - y) =$
(if $uint\ y \leq uint\ x$ *then*
 $uint\ x - uint\ y$
else
 $uint\ x - uint\ y + 2^{size\ x}$)
 ⟨*proof*⟩

lemma *unat-sub*:

$b \leq a \implies unat\ (a - b) = unat\ a - unat\ b$
 ⟨*proof*⟩

lemmas *word-less-sub1-numberof* [*simp*] = *word-less-sub1* [*of numeral w*] **for** w

lemmas *word-le-sub1-numberof* [*simp*] = *word-le-sub1* [*of numeral w*] **for** w

lemma *word-of-int-minus*:

$word-of-int\ (2^{len-of\ TYPE('a)} - i) = (word-of-int\ (-i)::'a::len\ word)$
 ⟨*proof*⟩

lemmas *word-of-int-inj* =

$word-uint.Abs-inject$ [*unfolded uints-num, simplified*]

lemma *word-le-less-eq*:

$(x::'z::len\ word) \leq y = (x = y \vee x < y)$

<proof>

lemma *mod-plus-cong*:

assumes 1: $(b::int) = b'$
and 2: $x \bmod b' = x' \bmod b'$
and 3: $y \bmod b' = y' \bmod b'$
and 4: $x' + y' = z'$
shows $(x + y) \bmod b = z' \bmod b'$
<proof>

lemma *mod-minus-cong*:

assumes 1: $(b::int) = b'$
and 2: $x \bmod b' = x' \bmod b'$
and 3: $y \bmod b' = y' \bmod b'$
and 4: $x' - y' = z'$
shows $(x - y) \bmod b = z' \bmod b'$
<proof>

lemma *word-induct-less*:

$\llbracket P (0::'a::len\ word); \bigwedge n. \llbracket n < m; P\ n \rrbracket \implies P\ (1 + n) \rrbracket \implies P\ m$
<proof>

lemma *word-induct*:

$\llbracket P (0::'a::len\ word); \bigwedge n. P\ n \implies P\ (1 + n) \rrbracket \implies P\ m$
<proof>

lemma *word-induct2* [*induct type*]:

$\llbracket P\ 0; \bigwedge n. \llbracket 1 + n \neq 0; P\ n \rrbracket \implies P\ (1 + n) \rrbracket \implies P\ (n::'b::len\ word)$
<proof>

16.29 Recursion combinator for words

definition *word-rec* :: $'a \Rightarrow ('b::len\ word \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'b\ word \Rightarrow 'a$
where

$word\text{-}rec\ forZero\ forSuc\ n = rec\text{-}nat\ forZero\ (forSuc \circ of\text{-}nat)\ (unat\ n)$

lemma *word-rec-0*: $word\text{-}rec\ z\ s\ 0 = z$

<proof>

lemma *word-rec-Suc*:

$1 + n \neq (0::'a::len\ word) \implies word\text{-}rec\ z\ s\ (1 + n) = s\ n\ (word\text{-}rec\ z\ s\ n)$
<proof>

lemma *word-rec-Pred*:

$n \neq 0 \implies word\text{-}rec\ z\ s\ n = s\ (n - 1)\ (word\text{-}rec\ z\ s\ (n - 1))$
<proof>

lemma *word-rec-in*:

$f\ (word\text{-}rec\ z\ (\lambda\cdot. f)\ n) = word\text{-}rec\ (f\ z)\ (\lambda\cdot. f)\ n$

<proof>

lemma *word-rec-in2*:

$f\ n\ (\text{word-rec}\ z\ f\ n) = \text{word-rec}\ (f\ 0\ z)\ (f\ \circ\ op\ +\ 1)\ n$
<proof>

lemma *word-rec-twice*:

$m \leq n \implies \text{word-rec}\ z\ f\ n = \text{word-rec}\ (\text{word-rec}\ z\ f\ (n - m))\ (f\ \circ\ op\ +\ (n - m))\ m$
<proof>

lemma *word-rec-id*: $\text{word-rec}\ z\ (\lambda\ \cdot.\ id)\ n = z$

<proof>

lemma *word-rec-id-eq*: $\forall m < n.\ f\ m = id \implies \text{word-rec}\ z\ f\ n = z$

<proof>

lemma *word-rec-max*:

$\forall m \geq n.\ m \neq -1 \implies f\ m = id \implies \text{word-rec}\ z\ f\ (-1) = \text{word-rec}\ z\ f\ n$
<proof>

lemma *unatSuc*:

$1 + n \neq (0::'a::\text{len}\ \text{word}) \implies \text{unat}\ (1 + n) = \text{Suc}\ (\text{unat}\ n)$
<proof>

declare *bin-to-bl-def* [*simp*]

<ML>

hide-const (**open**) *Word*

end

References

- [1] Jeremy Dawson. Isabelle theories for machine words. In Michael Goldsmith and Bill Roscoe, editors, *Seventh International Workshop on Automated Verification of Critical Systems (AVOCS'07)*, Electronic Notes in Theoretical Computer Science, page 15, Oxford, September 2007. Elsevier. to appear.