

# Transitive Union-Closed Families

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## Abstract

We formalise a proof by Aaronson, Ellis and Leader showing that the Union-Closed Conjecture holds for the union-closed family generated by the cyclic translates of any fixed set.

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## 1 Transitive Union-Closed Families

A family of sets is union-closed if the union of any two sets from the family is in the family. The Union-Closed Conjecture is an open problem in combinatorics posed by Frankl in 1979. It states that for every finite, union-closed family of sets (other than the family containing only the empty set) there exists an element that belongs to at least half of the sets in the family. We formalise a proof by Aaronson, Ellis and Leader showing that the Union-Closed Conjecture holds for the union-closed family generated by the cyclic translates of any fixed set [1].

**theory** *Transitive-Union-Closed-Families*

**imports** *Pluennecke-Ruzsa-Inequality.Pluennecke-Ruzsa-Inequality*

**begin**

**no-notation** *equivalence.Partition (infixl '/' 75)*

**definition** *union-closed:: 'a set set  $\Rightarrow$  bool*

**where** *union-closed  $\mathcal{F} \equiv (\forall A \in \mathcal{F}. \forall B \in \mathcal{F}. A \cup B \in \mathcal{F})$*

**abbreviation** *set-difference* :: [*'a set, 'a set*]  $\Rightarrow$  *'a set* (**infixl** \ 65)  
**where**  $A \setminus B \equiv A - B$

**locale** *Family* = *additive-abelian-group* +  
**fixes** *R*  
**assumes** *finG*: *finite G*  
**assumes** *RG*:  $R \subseteq G$   
**assumes** *R-nonempty*:  $R \neq \{\}$

**begin**

**definition** *union-closed-conjecture-property*:: *'a set set*  $\Rightarrow$  *bool*  
**where** *union-closed-conjecture-property*  $\mathcal{F}$   
 $\equiv \exists \mathcal{X} \subseteq \mathcal{F}. \exists x \in G. x \in \bigcap \mathcal{X} \wedge \text{card } \mathcal{X} \geq \text{card } \mathcal{F} / 2$

**definition** *Neighbd*  $\equiv \lambda A. \text{sumset } A \ R$

**definition** *Interior*  $\equiv \lambda A. \{x \in G. \text{sumset } \{x\} \ R \subseteq A\}$

**definition**  $\mathcal{F} \equiv \text{Neighbd } ' \text{Pow } G$

the family  $\mathcal{F}$  as defined above and appears in the statement of the theorem [1] is finite, nonempty union-closed family .

**lemma** *cardF-gt0* [*simp*]:  $\text{card } \mathcal{F} > 0$  **and** *finiteF*: *finite*  $\mathcal{F}$   
**using** *F-def finG* **by** *fastforce+*

As a remark, we note that  $\mathcal{F}$  is nontrivial.

**lemma**  $\mathcal{F} \neq \{\{\}\}$   
**unfolding** *F-def image-def Neighbd-def set-eq-iff*  
**apply** *simp*  
**by** (*metis RG R-nonempty Pow-top disjoint-iff emptyE subset-eq sumset-is-empty-iff*)

**lemma** *union-closed*  $\mathcal{F}$

**proof**–

**have**  $*\forall A \subseteq G. \forall B \subseteq G. (\text{sumset } A \ R) \cup (\text{sumset } B \ R) = \text{sumset } (A \cup B) \ R$   
**by** (*simp add: sumset-subset-Un1*)  
**show** *?thesis* **using**  $*$   
**by** (*auto simp: union-closed-def F-def Neighbd-def*)

**qed**

**lemma** *cardG-gt0*:  $\text{card } G > 0$   
**using** *RG R-nonempty card-0-eq finG* **by** *blast*

**lemma** *F-subset*:  $\mathcal{F} \subseteq \text{Pow } G$

**by** (*simp add: Neighbd-def PowI F-def image-subset-iff sumset-subset-carrier*)

## 1.1 Proof of the main theorem

**lemma** *card-Interior-le*:

**assumes**  $S \subseteq G$   
**shows**  $\text{card } (\text{Interior } S) \leq \text{card } S$   
**proof** –  
**obtain**  $r$  **where**  $r \in R$   
**using**  $R\text{-nonempty}$  **by**  $\text{blast}$   
**show**  $?thesis$   
**proof** ( $\text{intro card-inj-on-le}$ )  
**let**  $?f = (\lambda x. x \oplus r)$   
**show**  $\text{inj-on } ?f (\text{Interior } S) ?f ' \text{Interior } S \subseteq S$   
**using**  $RG \langle r \in R \rangle$  **by** ( $\text{auto simp: Interior-def inj-on-def}$ )  
**show**  $\text{finite } S$   
**using**  $\text{assms finG finite-subset}$  **by**  $\text{blast}$   
**qed**  
**qed**

**lemma**  $\text{Interior-subset-G [iff]: Interior } S \subseteq G$   
**using**  $\text{Interior-def}$  **by**  $\text{auto}$

**lemma**  $\text{Neighbd-subset-G [iff]: Neighbd } S \subseteq G$   
**by** ( $\text{simp add: Neighbd-def sumset-subset-carrier}$ )

**lemma**  $\text{average-ge:}$   
**shows**  $(\sum S \in \mathcal{F}. (\text{card } S)) / \text{card } \mathcal{F} \geq \text{card } G / 2$   
**proof** –

**define**  $f$  **where**  $f \equiv \lambda S. \text{minusset } (G \setminus \text{Interior } S)$

The following corresponds to (1) in the paper.

**have**  $1: \text{card } S + \text{card } (f S) \geq \text{card } G$  **if**  $S \subseteq G$  **for**  $S$

**proof** –

**have**  $\text{card } (f S) = \text{card } G - \text{card } (\text{Interior } S)$

**unfolding**  $f\text{-def}$

**by** ( $\text{metis Diff-subset Interior-subset-G card-Diff-subset card-minusset' finG finite-subset}$ )

**with that show**  $?thesis$  **using**  $\text{card-Interior-le}$

**by** ( $\text{metis (no-types, lifting) add.commute diff-le-mono2 le-diff-conv}$ )

**qed**

The following corresponds to (2) in the paper.

**have**  $2: f S = \text{sumset } (\text{minusset } (G \setminus S)) R$  **if**  $S \subseteq G$  **for**  $S$

**proof** –

**have**  $*$ :  $x \in f S \iff x \in \text{sumset } (\text{minusset } (G \setminus S)) R$  **if**  $x \in G$  **for**  $x$

**proof** –

**have**  $x \in f S \iff \text{inverse } x \notin \text{Interior } S$

**using**  $\text{that minusset.simps}$  **by** ( $\text{fastforce simp: f-def}$ ) $+$

**also have**  $\dots \iff (\text{sumset } \{\text{inverse } x\} R) \cap (G \setminus S) \neq \{\}$

**using**  $\text{sumset-subset-carrier that}$  **by** ( $\text{auto simp: Interior-def}$ )

**also have**  $\dots \iff x \in \text{sumset } (\text{minusset } (G \setminus S)) R$

**proof**

**assume**  $L: \text{sumset } \{\text{inverse } x\} R \cap (G \setminus S) \neq \{\}$

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then obtain  $r$  where  $r: \text{inverse } x \oplus r \notin S$  and  $r \in R$ 
  using  $\langle S \subseteq G \rangle \langle x \in G \rangle$  by (auto simp: sumset-eq minusset-eq)
then have  $\text{inverse } (\text{inverse } x \oplus r) \in \text{minusset } (G \setminus S)$ 
  using  $RG$  that by auto
moreover have  $x = \text{inverse } (\text{inverse } x \oplus r) \oplus r$ 
  using  $RG \langle r \in R \rangle$  that commutative inverse-composition-commute invert-
ible-right-inverse2
  by auto
ultimately show  $x \in \text{sumset } (\text{minusset } (G \setminus S)) R$ 
  by (metis  $RG \langle r \in R \rangle \text{minusset-subset-carrier subset-eq sumset.simps}$ )
next
assume  $R: x \in \text{sumset } (\text{minusset } (G \setminus S)) R$ 
then obtain  $g r$  where  $g \in G \ g \notin S \ r \in R \ x = \text{inverse } g \oplus r$ 
  by (metis  $\text{Diff-iff minusset.simps sumset.cases}$ )
show  $\text{sumset } \{\text{inverse } x\} R \cap (G \setminus S) \neq \{\}$ 
proof
  assume  $\text{sumset } \{\text{inverse } x\} R \cap (G \setminus S) = \{\}$ 
  then have  $g \notin \text{sumset } \{\text{inverse } x\} R$ 
    using  $\langle g \notin S \rangle \text{sumset-subset-carrier}$  that by fastforce
  then have  $g \neq \text{local.inverse } (\text{local.inverse } g \oplus r) \oplus r$ 
    using  $* RG$  that by (auto simp: sumset-eq)
  with  $* RG$  that show False
  by (metis  $\text{commutative invertible invertible-left-inverse2 invertible-right-inverse2}$ 
subset-eq)
  qed
qed
finally show ?thesis .
qed
show ?thesis
proof
  show  $f S \subseteq \text{sumset } (\text{minusset } (G \setminus S)) R$ 
    using  $* f\text{-def minusset-subset-carrier}$  by blast
next
  show  $\text{sumset } (\text{minusset } (G \setminus S)) R \subseteq f S$ 
    by (meson  $* \text{subset-iff sumset-subset-carrier}$ )
  qed
qed
then have  $f' \text{Pow } G \subseteq \mathcal{F}$ 
  by (auto simp:  $\text{Neighbd-def } \mathcal{F}\text{-def minusset-subset-carrier}$ )

  The following corresponds to (3) in the paper.
have  $3: \text{Neighbd } (\text{Interior } (\text{sumset } A R)) = \text{sumset } A R$ 
if  $A \subseteq G$  for  $A$ 
  using that by (force simp: sumset-eq Neighbd-def Interior-def)

  "Putting everything together":
moreover
have  $\text{sumset } X R = \text{sumset } Y R$ 
if  $X \subseteq G \ Y \subseteq G$ 

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$\text{minusset } (G \setminus \text{Interior } (\text{sumset } X R)) = \text{minusset } (G \setminus \text{Interior } (\text{sumset } Y R))$   
**for**  $X Y$   
**using** *that 3*  
**by** (*metis Diff-Diff-Int Int-absorb2 Interior-subset-G inf-commute minus-minusset*)  
**ultimately have**  $\text{inj-on } f \mathcal{F}$   
**by** (*auto simp: inj-on-def \mathcal{F}-def f-def Neighbd-def*)  
**moreover have**  $f' \mathcal{F} \subseteq \mathcal{F}$   
**using**  $2 \mathcal{F}\text{-def } \langle f' \text{ Pow } G \subseteq \mathcal{F} \rangle$  **by force**  
**moreover have**  $\mathcal{F} \subseteq f' \mathcal{F}$   
**by** (*metis <inj-on f \mathcal{F}> <f' \mathcal{F} \subseteq \mathcal{F}> endo-inj-surj finite\mathcal{F}*)  
**ultimately have**  $\text{bij-betw } f \mathcal{F} \mathcal{F}$   
**by** (*simp add: bij-betw-def*)  
**then have**  $\text{sum-card-eq: } (\sum S \in \mathcal{F}. \text{card } (f S)) = (\sum S \in \mathcal{F}. \text{card } S)$   
**by** (*simp add: sum.reindex-bij-betw*)  
  
**have**  $\text{card } G / 2 = (1 / (2 * \text{card } \mathcal{F})) * (\sum S \in \mathcal{F}. \text{card } G)$   
**by** *simp*  
**also have**  $\dots \leq (1 / (2 * \text{card } \mathcal{F})) * (\sum S \in \mathcal{F}. \text{card } S + \text{card } (f S))$   
**by** (*intro sum-mono mult-left-mono of-nat-mono 1*) (*auto simp: \mathcal{F}-def*)  
**also have**  $\dots = (1 / \text{card } \mathcal{F}) * (\sum S \in \mathcal{F}. \text{card } S)$   
**by** (*simp add: sum-card-eq sum.distrib*)  
**finally show** *?thesis*  
**by** *argo*  
**qed**

We have thus shown that the average size of a set in the family  $\mathcal{F}$  is at least  $|G|/2$ , proving the first part of Theorem 2 in the paper [1]. Using this, we will now show the main statement, i.e. that the Union-Closed Conjecture holds for the family  $\mathcal{F}$ .

**theorem** *Aaronson-Ellis-Leader-union-closed-conjecture:*

**shows** *union-closed-conjecture-property \mathcal{F}*

**proof** –

– First, quite a big calculation not mentioned in the article: counting all the elements in two different ways.

**have**  $*$ :  $(\sum S \in \mathcal{F}. (\text{card } S)) = (\sum x \in G. \text{card } \{S \in \mathcal{F}. x \in S\})$

**using** *finite\mathcal{F} \mathcal{F}-subset*

**proof** *induction*

**case** *empty*

**then show** *?case*

**by** *simp*

**next**

**case** (*insert S \mathcal{G}*)

**then have**  $A: \{T. (T = S \vee T \in \mathcal{G}) \wedge x \in T\}$

$= \{T \in \mathcal{G}. x \in T\} \cup (\text{if } x \in S \text{ then } \{S\} \text{ else } \{\})$

**for**  $x$

**by** *auto*

**have**  $B: \text{card } \{T. (T = S \vee T \in \mathcal{G}) \wedge x \in T\}$

$= \text{card } \{T \in \mathcal{G}. x \in T\} + (\text{if } x \in S \text{ then } 1 \text{ else } 0)$

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    for x
      by (simp add: A card-insert-if insert)
    have S = (∪ x∈G. if x ∈ S then {x} else {})
      using insert.premis by auto
    then have card S = card (∪ x∈G. if x ∈ S then {x} else {})
      by simp
    also have ... = (∑ i∈G. card (if i ∈ S then {i} else {}))
      by (intro card-UN-disjoint) (auto simp: finG)
    also have ... = (∑ x∈G. if x ∈ S then 1 else 0)
      by (force intro: sum.cong)
    finally have C: card S = (∑ x∈G. if x ∈ S then 1 else 0) .
    show ?case
      using insert by (auto simp: sum.distrib B C)
qed

have 1/2 ≤ (sum card F) / (card F * card G)
  using mult-right-mono [OF average-ge, of 1 / card G]
  using cardG-gt0 by (simp add: divide-simps split: if-splits)
also have ... = (∑ x∈G. ((card {S∈F. x∈S}) / (card F))) / card G
  by (simp add: * sum-divide-distrib)
finally have **: 1/2 ≤ (∑ x∈G. card {S∈F. x∈S} / card F) / card G .
  — There is a typo in the paper (bottom of page): instead of x ∈ S it says x ∈
  F.
show ?thesis
proof (rule ccontr) — Contradict the inequality proved above
  assume ¬ union-closed-conjecture-property F
  then have A: ∧ X x. [X ⊆ F; x ∈ G; x ∈ ∩ X] ⇒ card X < card F / 2
    by (fastforce simp: union-closed-conjecture-property-def)
  have (∑ x∈G. real (card {S∈F. x∈S})) < (∑ x∈G. card F / 2)
  proof (intro sum-strict-mono)
    fix x :: 'a
    assume x ∈ G
    then have card {S∈F. x∈S} < card F / 2
      by (intro A) auto
    then show real (card {S∈F. x∈S}) < real (card F) / 2
      by blast
  qed (use unit-closed finG in auto)
  also have ... = card F * (card G / 2)
    by simp
  finally have B: (∑ x∈G. real (card {S∈F. x∈S})) < card F * (card G / 2) .
  have (∑ x∈G. card {S∈F. x∈S} / card F) / card G < 1/2
    using cardG-gt0 divide-strict-right-mono [OF B, of card F * card G]
    by (simp add: divide-simps sum-divide-distrib)
  with ** show False
    by argo
  qed
qed
end

```

end

## References

- [1] J. Aaronson, D. Ellis, and I. Leader. A note on transitive union-closed families. 28(2), 2021. doi<https://doi.org/10.37236/9956>.