

Transitive Union-Closed Families

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February 1, 2025

Abstract

We formalise a proof by Aaronson, Ellis and Leader showing that the Union-Closed Conjecture holds for the union-closed family generated by the cyclic translates of any fixed set.

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1 Transitive Union-Closed Families

A family of sets is union-closed if the union of any two sets from the family is in the family. The Union-Closed Conjecture is an open problem in combinatorics posed by Frankl in 1979. It states that for every finite, union-closed family of sets (other than the family containing only the empty set) there exists an element that belongs to at least half of the sets in the family. We formalise a proof by Aaronson, Ellis and Leader showing that the Union-Closed Conjecture holds for the union-closed family generated by the cyclic translates of any fixed set [1].

theory *Transitive-Union-Closed-Families*

imports *Pluennecke-Ruzsa-Inequality.Pluennecke-Ruzsa-Inequality*

begin

no-notation *equivalence.Partition (infixl '/' 75)*

definition *union-closed:: 'a set set \Rightarrow bool*

where *union-closed $\mathcal{F} \equiv (\forall A \in \mathcal{F}. \forall B \in \mathcal{F}. A \cup B \in \mathcal{F})$*

abbreviation *set-difference* :: [*'a set, 'a set*] \Rightarrow *'a set* (**infixl** \ 65)
where $A \setminus B \equiv A - B$

locale *Family* = *additive-abelian-group* +
fixes *R*
assumes *finG*: *finite G*
assumes *RG*: $R \subseteq G$
assumes *R-nonempty*: $R \neq \{\}$

begin

definition *union-closed-conjecture-property*:: *'a set set* \Rightarrow *bool*
where *union-closed-conjecture-property* \mathcal{F}
 $\equiv \exists \mathcal{X} \subseteq \mathcal{F}. \exists x \in G. x \in \bigcap \mathcal{X} \wedge \text{card } \mathcal{X} \geq \text{card } \mathcal{F} / 2$

definition *Neighbd* $\equiv \lambda A. \text{sumset } A \ R$

definition *Interior* $\equiv \lambda A. \{x \in G. \text{sumset } \{x\} \ R \subseteq A\}$

definition $\mathcal{F} \equiv \text{Neighbd } ' \text{Pow } G$

We show that the family \mathcal{F} as defined above and appears in the statement of the theorem [1] is actually a finite, nonempty union-closed family indeed.

lemma *cardF-gt0* [*simp*]: $\text{card } \mathcal{F} > 0$ **and** *finiteF*: *finite* \mathcal{F}
using *F-def finG* **by** *fastforce+*

lemma *union-closed* \mathcal{F}

proof –

have $*\forall A \subseteq G. \forall B \subseteq G. (\text{sumset } A \ R) \cup (\text{sumset } B \ R) = \text{sumset } (A \cup B) \ R$
by (*simp add: sumset-subset-Un1*)

show *?thesis* **using** $*$

by (*auto simp: union-closed-def F-def Neighbd-def*)

qed

lemma *cardG-gt0*: $\text{card } G > 0$

using *RG R-nonempty card-0-eq finG* **by** *blast*

lemma *F-subset*: $\mathcal{F} \subseteq \text{Pow } G$

by (*simp add: Neighbd-def PowI F-def image-subset-iff sumset-subset-carrier*)

1.1 Proof of the main theorem

lemma *card-Interior-le*:

assumes $S \subseteq G$

shows $\text{card } (\text{Interior } S) \leq \text{card } S$

proof –

obtain r **where** $r \in R$

using *R-nonempty* **by** *blast*

show *?thesis*

```

proof (intro card-inj-on-le)
  let ?f = ( $\lambda x. x \oplus r$ )
  show inj-on ?f (Interior S) ?f ‘ Interior S  $\subseteq$  S
    using RG  $\langle r \in R \rangle$  by (auto simp: Interior-def inj-on-def)
  show finite S
    using assms finG finite-subset by blast
qed
qed

```

```

lemma Interior-subset-G [iff]: Interior S  $\subseteq$  G
  using Interior-def by auto

```

```

lemma Neighbd-subset-G [iff]: Neighbd S  $\subseteq$  G
  by (simp add: Neighbd-def sumset-subset-carrier)

```

```

lemma average-ge:

```

```

  shows ( $\sum S \in \mathcal{F}. (\text{card } S)$ ) / card  $\mathcal{F} \geq$  card G / 2

```

```

proof –

```

```

  define f where f  $\equiv \lambda S. \text{minusset } (G \setminus \text{Interior } S)$ 

```

The following corresponds to (1) in the paper.

```

  have 1: card S + card (f S)  $\geq$  card G if S  $\subseteq$  G for S

```

```

proof –

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```

  have card (f S) = card G – card (Interior S)

```

```

    unfolding f-def

```

```

    by (metis Diff-subset Interior-subset-G card-Diff-subset card-minusset' finG
  finite-subset)

```

```

    with that show ?thesis using card-Interior-le

```

```

    by (metis (no-types, lifting) add.commute diff-le-mono2 le-diff-conv)

```

```

qed

```

The following corresponds to (2) in the paper.

```

  have 2: f S = sumset (minusset (G \ S)) R if S  $\subseteq$  G for S

```

```

proof –

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```

  have *: x  $\in$  f S  $\iff$  x  $\in$  sumset (minusset (G \ S)) R if x  $\in$  G for x

```

```

proof –

```

```

  have x  $\in$  f S  $\iff$  inverse x  $\notin$  Interior S

```

```

    using that minusset.simps by (fastforce simp: f-def)+

```

```

  also have ...  $\iff$  (sumset {inverse x} R)  $\cap$  (G \ S)  $\neq$  {}

```

```

    using sumset-subset-carrier that by (auto simp: Interior-def)

```

```

  also have ...  $\iff$  x  $\in$  sumset (minusset (G \ S)) R

```

```

proof

```

```

  assume L: sumset {inverse x} R  $\cap$  (G \ S)  $\neq$  {}

```

```

  then obtain r where r: inverse x  $\oplus$  r  $\notin$  S and r  $\in$  R

```

```

    using  $\langle S \subseteq G \rangle \langle x \in G \rangle$  by (auto simp: sumset-eq minusset-eq)

```

```

  then have inverse (inverse x  $\oplus$  r)  $\in$  minusset (G \ S)

```

```

    using RG that by auto

```

```

  moreover have x = inverse (inverse x  $\oplus$  r)  $\oplus$  r

```

```

    using RG  $\langle r \in R \rangle$  that commutative inverse-composition-commute invert-
  ible-right-inverse2

```

```

    by auto
  ultimately show  $x \in \text{sumset } (\text{minusset } (G \setminus S)) R$ 
    by (metis  $RG \langle r \in R \rangle \text{minusset-subset-carrier subset-eq sumset.simps}$ )
next
  assume  $R: x \in \text{sumset } (\text{minusset } (G \setminus S)) R$ 
  then obtain  $g r$  where  $*: g \in G \ g \notin S \ r \in R \ x = \text{inverse } g \oplus r$ 
    by (metis  $\text{Diff-iff minusset.simps sumset.cases}$ )
  show  $\text{sumset } \{\text{inverse } x\} R \cap (G \setminus S) \neq \{\}$ 
  proof
    assume  $\text{sumset } \{\text{inverse } x\} R \cap (G \setminus S) = \{\}$ 
    then have  $g \notin \text{sumset } \{\text{inverse } x\} R$ 
      using  $\langle g \notin S \rangle \text{sumset-subset-carrier}$  that by fastforce
    then have  $g \neq \text{local.inverse } (\text{local.inverse } g \oplus r) \oplus r$ 
      using  $* RG$  that by (auto simp:  $\text{sumset-eq}$ )
    with  $* RG$  that show False
  by (metis  $\text{commutative invertible invertible-left-inverse2 invertible-right-inverse2}$ 
subset-eq)
  qed
  qed
  finally show ?thesis .
  qed
  show ?thesis
  proof
    show  $f S \subseteq \text{sumset } (\text{minusset } (G \setminus S)) R$ 
      using  $* f\text{-def minusset-subset-carrier}$  by blast
  next
    show  $\text{sumset } (\text{minusset } (G \setminus S)) R \subseteq f S$ 
      by (meson  $* \text{subset-iff sumset-subset-carrier}$ )
  qed
  qed
  then have  $f \text{ ` Pow } G \subseteq \mathcal{F}$ 
    by (auto simp:  $\text{Neighbd-def } \mathcal{F}\text{-def minusset-subset-carrier}$ )

  The following corresponds to (3) in the paper.

  have 3:  $\text{Neighbd } (\text{Interior } (\text{sumset } A R)) = \text{sumset } A R$ 
    if  $A \subseteq G$  for  $A$ 
    using that by (force simp:  $\text{sumset-eq Neighbd-def Interior-def}$ )

  "Putting everything together":

  moreover
  have  $\text{sumset } X R = \text{sumset } Y R$ 
    if  $X \subseteq G \ Y \subseteq G$ 
       $\text{minusset } (G \setminus \text{Interior } (\text{sumset } X R)) = \text{minusset } (G \setminus \text{Interior } (\text{sumset } Y$ 
 $R))$ 
    for  $X Y$ 
    using that 3
  by (metis  $\text{Diff-Diff-Int Int-absorb2 Interior-subset-G inf-commute minus-minusset}$ )
  ultimately have  $\text{inj-on } f \mathcal{F}$ 
    by (auto simp:  $\text{inj-on-def } \mathcal{F}\text{-def } f\text{-def Neighbd-def}$ )

```

moreover have $f' \mathcal{F} \subseteq \mathcal{F}$
using \mathcal{F} -def $\langle f' \text{Pow } G \subseteq \mathcal{F} \rangle$ **by force**
moreover have $\mathcal{F} \subseteq f' \mathcal{F}$
by (*metis* $\langle \text{inj-on } f \mathcal{F} \rangle \langle f' \mathcal{F} \subseteq \mathcal{F} \rangle$ *endo-inj-surj finite* \mathcal{F})
ultimately have *bij-betw* $f \mathcal{F} \mathcal{F}$
by (*simp add: bij-betw-def*)
then have *sum-card-eq*: $(\sum_{S \in \mathcal{F}} \text{card } (f S)) = (\sum_{S \in \mathcal{F}} \text{card } S)$
by (*simp add: sum.reindex-bij-betw*)

have $\text{card } G / 2 = (1 / (2 * \text{card } \mathcal{F})) * (\sum_{S \in \mathcal{F}} \text{card } G)$
by *simp*
also have $\dots \leq (1 / (2 * \text{card } \mathcal{F})) * (\sum_{S \in \mathcal{F}} \text{card } S + \text{card } (f S))$
by (*intro sum-mono mult-left-mono of-nat-mono 1*) (*auto simp: \mathcal{F}*-def)
also have $\dots = (1 / \text{card } \mathcal{F}) * (\sum_{S \in \mathcal{F}} \text{card } S)$
by (*simp add: sum-card-eq sum.distrib*)
finally show *?thesis*
by *argo*
qed

We have thus shown that the average size of a set in the family \mathcal{F} is at least $|G|/2$, proving the first part of Theorem 2 in the paper [1]. Using this, we will now show the main statement, i.e. that the Union-Closed Conjecture holds for the family \mathcal{F} .

theorem *Aaronson-Ellis-Leader-union-closed-conjecture*:

shows *union-closed-conjecture-property* \mathcal{F}

proof –

– First, quite a big calculation not mentioned in the article: counting all the elements in two different ways.

have $*$: $(\sum_{S \in \mathcal{F}} (\text{card } S)) = (\sum_{x \in G} \text{card } \{S \in \mathcal{F}. x \in S\})$

using *finite* \mathcal{F} *\mathcal{F}*-subset

proof *induction*

case *empty*

then show *?case*

by *simp*

next

case (*insert* $S \mathcal{G}$)

then have A : $\{T. (T = S \vee T \in \mathcal{G}) \wedge x \in T\}$
 $= \{T \in \mathcal{G}. x \in T\} \cup (\text{if } x \in S \text{ then } \{S\} \text{ else } \{\})$

for x

by *auto*

have B : $\text{card } \{T. (T = S \vee T \in \mathcal{G}) \wedge x \in T\}$
 $= \text{card } \{T \in \mathcal{G}. x \in T\} + (\text{if } x \in S \text{ then } 1 \text{ else } 0)$

for x

by (*simp add: A card-insert-if insert*)

have $S = (\bigcup_{x \in G} \text{if } x \in S \text{ then } \{x\} \text{ else } \{\})$

using *insert.premis* **by** *auto*

then have $\text{card } S = \text{card } (\bigcup_{x \in G} \text{if } x \in S \text{ then } \{x\} \text{ else } \{\})$

by *simp*

also have $\dots = (\sum_{i \in G} \text{card } (\text{if } i \in S \text{ then } \{i\} \text{ else } \{\}))$

```

    by (intro card-UN-disjoint) (auto simp: finG)
  also have ... = (∑ x∈G. if x ∈ S then 1 else 0)
    by (force intro: sum.cong)
  finally have C: card S = (∑ x∈G. if x ∈ S then 1 else 0) .
  show ?case
    using insert by (auto simp: sum.distrib B C)
qed

have 1/2 ≤ (sum card F) / (card F * card G)
  using mult-right-mono [OF average-ge, of 1 / card G]
  using cardG-gt0 by (simp add: divide-simps split: if-splits)
also have ... = (∑ x∈G. ((card {S∈F. x∈S}) / (card F))) / card G
  by (simp add: * sum-divide-distrib)
finally have **: 1/2 ≤ (∑ x∈G. card {S∈F. x∈S} / card F) / card G .
  — There is a typo in the paper (bottom of page): instead of  $x \in S$  it says  $x \in \mathcal{F}$ .
show ?thesis
proof (rule ccontr) — Contradict the inequality proved above
  assume ¬ union-closed-conjecture-property F
  then have A:  $\bigwedge \mathcal{X} x. [\mathcal{X} \subseteq \mathcal{F}; x \in G; x \in \bigcap \mathcal{X}] \implies \text{card } \mathcal{X} < \text{card } \mathcal{F} / 2$ 
    by (fastforce simp: union-closed-conjecture-property-def)
  have (∑ x∈G. real (card {S∈F. x∈S})) < (∑ x∈G. card F / 2)
  proof (intro sum-strict-mono)
    fix x :: 'a
    assume x ∈ G
    then have card {S∈F. x∈S} < card F / 2
      by (intro A) auto
    then show real (card {S∈F. x∈S}) < real (card F) / 2
      by blast
  qed (use unit-closed finG in auto)
  also have ... = card F * (card G / 2)
    by simp
  finally have B: (∑ x∈G. real (card {S∈F. x∈S})) < card F * (card G / 2) .
  have (∑ x∈G. card {S∈F. x∈S} / card F) / card G < 1/2
    using divide-strict-right-mono [OF B, of card F * card G]
    using cardG-gt0
    by (simp add: divide-simps sum-divide-distrib)
  with ** show False
    by argo
qed
qed
end
end

```

References

- [1] J. Aaronson, D. Ellis, and I. Leader. A note on transitive union-closed families. 28(2), 2021. doi<https://doi.org/10.37236/9956>.