

A Formalisation of Sturm's Theorem

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Abstract

Sturm sequences are a method for computing the number of real roots of a real polynomial inside a given interval efficiently. In this project, this fact and a number of methods to construct Sturm sequences efficiently have been formalised with the interactive theorem prover Isabelle/HOL. Building upon this, an Isabelle/HOL proof method was then implemented to prove statements about the number of roots of a real polynomial and related properties.

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1 Miscellaneous

```
theory Misc-Polynomial
imports HOL-Computational-Algebra.Polynomial HOL-Computational-Algebra.Polynomial-Factorial
Pure-ex.Guess
begin
```

1.1 Analysis

```
lemma fun-eq-in-ivl:
  assumes  $a \leq b \ \forall x::\text{real}. a \leq x \wedge x \leq b \longrightarrow \text{eventually } (\lambda \xi. f \ \xi = f \ x) \text{ (at } x)$ 
  shows  $f \ a = f \ b$ 
<proof>
```

1.2 Polynomials

1.2.1 General simplification lemmas

```
lemma pderiv-div:
  assumes [simp]:  $q \ \text{dvd} \ p \ \ q \neq 0$ 
  shows  $\text{pderiv } (p \ \text{div} \ q) = (q * \text{pderiv } p - p * \text{pderiv } q) \ \text{div} \ (q * q)$ 
 $q * q \ \text{dvd} \ (q * \text{pderiv } p - p * \text{pderiv } q)$ 
<proof>
```

1.2.2 Divisibility of polynomials

Two polynomials that are coprime have no common roots.

```
lemma coprime-imp-no-common-roots:
   $\neg (\text{poly } p \ x = 0 \wedge \text{poly } q \ x = 0) \ \text{if } \text{coprime } p \ q$ 
  for  $x :: 'a :: \text{field}$ 
<proof>
```

```
lemma poly-div:
  assumes  $\text{poly } q \ x \neq 0$  and  $(q::'a :: \text{field } \text{poly}) \ \text{dvd} \ p$ 
  shows  $\text{poly } (p \ \text{div} \ q) \ x = \text{poly } p \ x / \text{poly } q \ x$ 
<proof>
```

```
lemma poly-div-gcd-squarefree-aux:
  assumes  $\text{pderiv } (p::('a::\{\text{field-char-0,field-gcd}\}) \ \text{poly}) \neq 0$ 
  defines  $d \equiv \text{gcd } p \ (\text{pderiv } p)$ 
  shows  $\text{coprime } (p \ \text{div} \ d) \ (\text{pderiv } (p \ \text{div} \ d))$  and
 $\bigwedge x. \text{poly } (p \ \text{div} \ d) \ x = 0 \longleftrightarrow \text{poly } p \ x = 0$ 
<proof>
```

```
lemma normalize-field:
   $\text{normalize } (x :: 'a :: \{\text{field,normalization-semidom}\}) = (\text{if } x = 0 \text{ then } 0 \text{ else } 1)$ 
<proof>
```

lemma *normalize-field-eq-1* [simp]:

$x \neq 0 \implies \text{normalize } (x :: 'a :: \{\text{field}, \text{normalization-semidom}\}) = 1$
(proof)

lemma *unit-factor-field* [simp]:

$\text{unit-factor } (x :: 'a :: \{\text{field}, \text{normalization-semidom}\}) = x$
(proof)

Dividing a polynomial by its gcd with its derivative yields a squarefree polynomial with the same roots.

lemma *poly-div-gcd-squarefree*:

assumes $(p :: ('a :: \{\text{field-char-0}, \text{field-gcd}\}) \text{poly}) \neq 0$
defines $d \equiv \text{gcd } p \ (\text{pderiv } p)$
shows $\text{coprime } (p \ \text{div } d) \ (\text{pderiv } (p \ \text{div } d))$ (is ?A) **and**
 $\bigwedge x. \text{poly } (p \ \text{div } d) \ x = 0 \iff \text{poly } p \ x = 0$ (is $\bigwedge x. ?B \ x$)
(proof)

1.2.3 Sign changes of a polynomial

If a polynomial has different signs at two points, it has a root inbetween.

lemma *poly-different-sign-imp-root*:

assumes $a < b$ **and** $\text{sgn } (\text{poly } p \ a) \neq \text{sgn } (\text{poly } p \ (b :: \text{real}))$
shows $\exists x. a \leq x \wedge x \leq b \wedge \text{poly } p \ x = 0$
(proof)

lemma *poly-different-sign-imp-root'*:

assumes $\text{sgn } (\text{poly } p \ a) \neq \text{sgn } (\text{poly } p \ (b :: \text{real}))$
shows $\exists x. \text{poly } p \ x = 0$
(proof)

lemma *no-roots-inbetween-imp-same-sign*:

assumes $a < b \ \forall x. a \leq x \wedge x \leq b \implies \text{poly } p \ x \neq (0 :: \text{real})$
shows $\text{sgn } (\text{poly } p \ a) = \text{sgn } (\text{poly } p \ b)$
(proof)

1.2.4 Limits of polynomials

lemma *poly-neighbourhood-without-roots*:

assumes $(p :: \text{real poly}) \neq 0$
shows *eventually* $(\lambda x. \text{poly } p \ x \neq 0)$ (at x_0)
(proof)

lemma *poly-neighbourhood-same-sign*:

assumes $\text{poly } p \ (x_0 :: \text{real}) \neq 0$
shows *eventually* $(\lambda x. \text{sgn } (\text{poly } p \ x) = \text{sgn } (\text{poly } p \ x_0))$ (at x_0)
(proof)

lemma *poly-lhopital*:

assumes $\text{poly } p (x::\text{real}) = 0$ $\text{poly } q x = 0$ $q \neq 0$

assumes $(\lambda x. \text{poly } (\text{pderiv } p) x / \text{poly } (\text{pderiv } q) x) -x \rightarrow y$

shows $(\lambda x. \text{poly } p x / \text{poly } q x) -x \rightarrow y$

<proof>

lemma *poly-roots-bounds*:

assumes $p \neq 0$

obtains $l u$

where $l \leq (u :: \text{real})$

and $\text{poly } p l \neq 0$

and $\text{poly } p u \neq 0$

and $\{x. x > l \wedge x \leq u \wedge \text{poly } p x = 0\} = \{x. \text{poly } p x = 0\}$

and $\bigwedge x. x \leq l \implies \text{sgn } (\text{poly } p x) = \text{sgn } (\text{poly } p l)$

and $\bigwedge x. x \geq u \implies \text{sgn } (\text{poly } p x) = \text{sgn } (\text{poly } p u)$

<proof>

definition *poly-inf* :: $('a::\text{real-normed-vector}) \text{poly} \Rightarrow 'a$ **where**

$\text{poly-inf } p \equiv \text{sgn } (\text{coeff } p (\text{degree } p))$

definition *poly-neg-inf* :: $('a::\text{real-normed-vector}) \text{poly} \Rightarrow 'a$ **where**

$\text{poly-neg-inf } p \equiv \text{if even } (\text{degree } p) \text{ then } \text{sgn } (\text{coeff } p (\text{degree } p))$
 $\text{else } -\text{sgn } (\text{coeff } p (\text{degree } p))$

lemma *poly-inf-0-iff*[simp]:

$\text{poly-inf } p = 0 \iff p = 0$ $\text{poly-neg-inf } p = 0 \iff p = 0$

<proof>

lemma *poly-inf-mult*[simp]:

fixes $p :: ('a::\text{real-normed-field}) \text{poly}$

shows $\text{poly-inf } (p*q) = \text{poly-inf } p * \text{poly-inf } q$

$\text{poly-neg-inf } (p*q) = \text{poly-neg-inf } p * \text{poly-neg-inf } q$

<proof>

lemma *poly-neq-0-at-infinity*:

assumes $(p :: \text{real poly}) \neq 0$

shows *eventually* $(\lambda x. \text{poly } p x \neq 0)$ *at-infinity*

<proof>

lemma *poly-limit-aux*:

fixes $p :: \text{real poly}$

defines $n \equiv \text{degree } p$

shows $((\lambda x. \text{poly } p x / x^n) \longrightarrow \text{coeff } p n)$ *at-infinity*

<proof>

lemma *poly-at-top-at-top*:

fixes $p :: \text{real poly}$
assumes $\text{degree } p \geq 1 \text{ coeff } p (\text{degree } p) > 0$
shows $\text{LIM } x \text{ at-top. } \text{poly } p \ x \text{ :> at-top}$
(proof)

lemma *poly-at-bot-at-top*:

fixes $p :: \text{real poly}$
assumes $\text{degree } p \geq 1 \text{ coeff } p (\text{degree } p) < 0$
shows $\text{LIM } x \text{ at-top. } \text{poly } p \ x \text{ :> at-bot}$
(proof)

lemma *poly-lim-inf*:

eventually $(\lambda x :: \text{real. } \text{sgn } (\text{poly } p \ x) = \text{poly-inf } p) \text{ at-top}$
(proof)

lemma *poly-at-top-or-bot-at-bot*:

fixes $p :: \text{real poly}$
assumes $\text{degree } p \geq 1 \text{ coeff } p (\text{degree } p) > 0$
shows $\text{LIM } x \text{ at-bot. } \text{poly } p \ x \text{ :> (if even (degree } p) \text{ then at-top else at-bot)}$
(proof)

lemma *poly-at-bot-or-top-at-bot*:

fixes $p :: \text{real poly}$
assumes $\text{degree } p \geq 1 \text{ coeff } p (\text{degree } p) < 0$
shows $\text{LIM } x \text{ at-bot. } \text{poly } p \ x \text{ :> (if even (degree } p) \text{ then at-bot else at-top)}$
(proof)

lemma *poly-lim-neg-inf*:

eventually $(\lambda x :: \text{real. } \text{sgn } (\text{poly } p \ x) = \text{poly-neg-inf } p) \text{ at-bot}$
(proof)

1.2.5 Signs of polynomials for sufficiently large values

lemma *polys-inf-sign-thresholds*:

assumes *finite* ($ps :: \text{real poly set}$)
obtains $l \ u$
where $l \leq u$
and $\bigwedge p. \llbracket p \in ps; p \neq 0 \rrbracket \implies$
 $\{x. l < x \wedge x \leq u \wedge \text{poly } p \ x = 0\} = \{x. \text{poly } p \ x = 0\}$
and $\bigwedge p \ x. \llbracket p \in ps; x \geq u \rrbracket \implies \text{sgn } (\text{poly } p \ x) = \text{poly-inf } p$
and $\bigwedge p \ x. \llbracket p \in ps; x \leq l \rrbracket \implies \text{sgn } (\text{poly } p \ x) = \text{poly-neg-inf } p$
(proof)

1.2.6 Positivity of polynomials

lemma *poly-pos*:

$$(\forall x::real. poly\ p\ x > 0) \longleftrightarrow poly\text{-}inf\ p = 1 \wedge (\forall x. poly\ p\ x \neq 0)$$

<proof>

lemma *poly-pos-greater*:

$$(\forall x::real. x > a \longrightarrow poly\ p\ x > 0) \longleftrightarrow$$
$$poly\text{-}inf\ p = 1 \wedge (\forall x. x > a \longrightarrow poly\ p\ x \neq 0)$$

<proof>

lemma *poly-pos-geq*:

$$(\forall x::real. x \geq a \longrightarrow poly\ p\ x > 0) \longleftrightarrow$$
$$poly\text{-}inf\ p = 1 \wedge (\forall x. x \geq a \longrightarrow poly\ p\ x \neq 0)$$

<proof>

lemma *poly-pos-less*:

$$(\forall x::real. x < a \longrightarrow poly\ p\ x > 0) \longleftrightarrow$$
$$poly\text{-}neg\text{-}inf\ p = 1 \wedge (\forall x. x < a \longrightarrow poly\ p\ x \neq 0)$$

<proof>

lemma *poly-pos-leq*:

$$(\forall x::real. x \leq a \longrightarrow poly\ p\ x > 0) \longleftrightarrow$$
$$poly\text{-}neg\text{-}inf\ p = 1 \wedge (\forall x. x \leq a \longrightarrow poly\ p\ x \neq 0)$$

<proof>

lemma *poly-pos-between-less-less*:

$$(\forall x::real. a < x \wedge x < b \longrightarrow poly\ p\ x > 0) \longleftrightarrow$$
$$(a \geq b \vee poly\ p\ ((a+b)/2) > 0) \wedge (\forall x. a < x \wedge x < b \longrightarrow poly\ p\ x \neq 0)$$

<proof>

lemma *poly-pos-between-less-leq*:

$$(\forall x::real. a < x \wedge x \leq b \longrightarrow poly\ p\ x > 0) \longleftrightarrow$$
$$(a \geq b \vee poly\ p\ b > 0) \wedge (\forall x. a < x \wedge x \leq b \longrightarrow poly\ p\ x \neq 0)$$

<proof>

lemma *poly-pos-between-leq-less*:

$$(\forall x::real. a \leq x \wedge x < b \longrightarrow poly\ p\ x > 0) \longleftrightarrow$$
$$(a \geq b \vee poly\ p\ a > 0) \wedge (\forall x. a \leq x \wedge x < b \longrightarrow poly\ p\ x \neq 0)$$

<proof>

lemma *poly-pos-between-leq-leq*:

$$(\forall x::real. a \leq x \wedge x \leq b \longrightarrow poly\ p\ x > 0) \longleftrightarrow$$
$$(a > b \vee poly\ p\ a > 0) \wedge (\forall x. a \leq x \wedge x \leq b \longrightarrow poly\ p\ x \neq 0)$$

<proof>

end

2 Proof of Sturm's Theorem

```

theory Sturm-Theorem
  imports HOL-Computational-Algebra.Polynomial
           Lib/Sturm-Library HOL-Computational-Algebra.Field-as-Ring
begin

```

2.1 Sign changes of polynomial sequences

For a given sequence of polynomials, this function computes the number of sign changes of the sequence of polynomials evaluated at a given position x . A sign change is a change from a negative value to a positive one or vice versa; zeros in the sequence are ignored.

definition *sign-changes where*

```

sign-changes ps (x::real) =
  length (remdups-adj (filter (λx. x ≠ 0) (map (λp. sgn (poly p x)) ps))) - 1

```

The number of sign changes of a sequence distributes over a list in the sense that the number of sign changes of a sequence $p_1, \dots, p_i, \dots, p_n$ at x is the same as the sum of the sign changes of the sequence p_1, \dots, p_i and p_i, \dots, p_n as long as $p_i(x) \neq 0$.

lemma *sign-changes-distrib:*

```

poly p x ≠ 0 ⇒
  sign-changes (ps1 @ [p] @ ps2) x =
  sign-changes (ps1 @ [p]) x + sign-changes ([p] @ ps2) x
⟨proof⟩

```

The following two congruences state that the number of sign changes is the same if all the involved signs are the same.

lemma *sign-changes-cong:*

```

assumes length ps = length ps'
assumes ∀ i < length ps. sgn (poly (ps!i) x) = sgn (poly (ps'!i) y)
shows sign-changes ps x = sign-changes ps' y
⟨proof⟩

```

lemma *sign-changes-cong':*

```

assumes ∀ p ∈ set ps. sgn (poly p x) = sgn (poly p y)
shows sign-changes ps x = sign-changes ps y
⟨proof⟩

```

For a sequence of polynomials of length 3, if the first and the third polynomial have opposite and nonzero sign at some x , the number of sign changes is always 1, irrespective of the sign of the second polynomial.

lemma *sign-changes-sturm-triple:*

```

assumes poly p x ≠ 0 and sgn (poly r x) = - sgn (poly p x)
shows sign-changes [p,q,r] x = 1
⟨proof⟩

```

Finally, we define two additional functions that count the sign changes “at infinity”.

definition *sign-changes-inf* **where**

$$\text{sign-changes-inf } ps = \text{length } (\text{remdups-adj } (\text{filter } (\lambda x. x \neq 0) (\text{map } \text{poly-inf } ps))) - 1$$

definition *sign-changes-neg-inf* **where**

$$\text{sign-changes-neg-inf } ps = \text{length } (\text{remdups-adj } (\text{filter } (\lambda x. x \neq 0) (\text{map } \text{poly-neg-inf } ps))) - 1$$

2.2 Definition of Sturm sequences locale

We first define the notion of a “Quasi-Sturm sequence”, which is a weakening of a Sturm sequence that captures the properties that are fulfilled by a nonempty suffix of a Sturm sequence:

- The sequence is nonempty.
- The last polynomial does not change its sign.
- If the middle one of three adjacent polynomials has a root at x , the other two have opposite and nonzero signs at x .

locale *quasi-sturm-seq* =

fixes $ps :: (\text{real poly}) \text{ list}$

assumes $\text{last-ps-sgn-const}[simp]$:

$$\bigwedge x y. \text{sgn } (\text{poly } (\text{last } ps) x) = \text{sgn } (\text{poly } (\text{last } ps) y)$$

assumes $\text{ps-not-Nil}[simp]$: $ps \neq []$

assumes signs : $\bigwedge i x. \llbracket i < \text{length } ps - 2; \text{poly } (ps ! (i+1)) x = 0 \rrbracket$
 $\implies (\text{poly } (ps ! (i+2)) x) * (\text{poly } (ps ! i) x) < 0$

Now we define a Sturm sequence p_1, \dots, p_n of a polynomial p in the following way:

- The sequence contains at least two elements.
- p is the first polynomial, i. e. $p_1 = p$.
- At any root x of p , p_2 and p have opposite sign left of x and the same sign right of x in some neighbourhood around x .
- The first two polynomials in the sequence have no common roots.
- If the middle one of three adjacent polynomials has a root at x , the other two have opposite and nonzero signs at x .

locale *sturm-seq* = *quasi-sturm-seq* +

fixes $p :: \text{real poly}$

```

assumes hd-ps-p[simp]: hd ps = p
assumes length-ps-ge-2[simp]: length ps ≥ 2
assumes deriv:  $\bigwedge x_0. \text{poly } p \ x_0 = 0 \implies$ 
    eventually  $(\lambda x. \text{sgn } (\text{poly } (p * \text{ps!1}) \ x) =$ 
     $(\text{if } x > x_0 \ \text{then } 1 \ \text{else } -1)) \ (\text{at } x_0)$ 
assumes p-squarefree:  $\bigwedge x. \neg(\text{poly } p \ x = 0 \wedge \text{poly } (\text{ps!1}) \ x = 0)$ 
begin

```

Any Sturm sequence is obviously a Quasi-Sturm sequence.

```

lemma quasi-sturm-seq: quasi-sturm-seq ps <proof><proof><proof><proof>end
<proof>

```

Any suffix of a Quasi-Sturm sequence is again a Quasi-Sturm sequence.

```

lemma quasi-sturm-seq-Cons:
  assumes quasi-sturm-seq (p#ps) and ps ≠ []
  shows quasi-sturm-seq ps
<proof>

```

2.3 Auxiliary lemmas about roots and sign changes

```

lemma sturm-adjacent-root-aux:
  assumes i < length (ps :: real poly list) - 1
  assumes poly (ps ! i) x = 0 and poly (ps ! (i + 1)) x = 0
  assumes  $\bigwedge i \ x. \llbracket i < \text{length } ps - 2; \text{poly } (ps ! (i+1)) \ x = 0 \rrbracket$ 
     $\implies \text{sgn } (\text{poly } (ps ! (i+2)) \ x) = - \text{sgn } (\text{poly } (ps ! i) \ x)$ 
  shows  $\forall j \leq i+1. \text{poly } (ps ! j) \ x = 0$ 
<proof>

```

This function splits the sign list of a Sturm sequence at a position x that is not a root of p into a list of sublists such that the number of sign changes within every sublist is constant in the neighbourhood of x , thus proving that the total number is also constant.

```

fun split-sign-changes where
  split-sign-changes [p] (x :: real) = [[p]] |
  split-sign-changes [p,q] x = [[p,q]] |
  split-sign-changes (p#q#r#ps) x =
    (if poly p x ≠ 0 ∧ poly q x = 0 then
      [p,q,r] # split-sign-changes (r#ps) x
    else
      [p,q] # split-sign-changes (q#r#ps) x)

```

```

lemma (in quasi-sturm-seq) split-sign-changes-subset[dest]:
  ps' ∈ set (split-sign-changes ps x) ⟹ set ps' ⊆ set ps
<proof>

```

A custom induction rule for *split-sign-changes* that uses the fact that all the intermediate parameters in calls of *split-sign-changes* are quasi-Sturm sequences.

lemma (in *quasi-sturm-seq*) *split-sign-changes-induct*:

$$\llbracket \bigwedge p x. P [p] x; \bigwedge p q x. \text{quasi-sturm-seq} [p, q] \implies P [p, q] x; \bigwedge p q r ps x. \text{quasi-sturm-seq} (p \# q \# r \# ps) \implies \llbracket \text{poly } p x \neq 0 \implies \text{poly } q x = 0 \implies P (r \# ps) x; \text{poly } q x \neq 0 \implies P (q \# r \# ps) x; \text{poly } p x = 0 \implies P (q \# r \# ps) x \rrbracket \implies P (p \# q \# r \# ps) x \rrbracket \implies P ps x$$

<proof>

The total number of sign changes in the split list is the same as the number of sign changes in the original list.

lemma (in *quasi-sturm-seq*) *split-sign-changes-correct*:
assumes $\text{poly } (hd \ ps) \ x_0 \neq 0$
defines $\text{sign-changes}' \equiv \lambda ps \ x. \sum ps' \leftarrow \text{split-sign-changes } ps \ x. \text{sign-changes } ps' \ x$
shows $\text{sign-changes}' \ ps \ x_0 = \text{sign-changes } ps \ x_0$

<proof>

We now prove that if $p(x) \neq 0$, the number of sign changes of a Sturm sequence of p at x is constant in a neighbourhood of x .

lemma (in *quasi-sturm-seq*) *split-sign-changes-correct-nbh*:
assumes $\text{poly } (hd \ ps) \ x_0 \neq 0$
defines $\text{sign-changes}' \equiv \lambda x_0 \ ps \ x. \sum ps' \leftarrow \text{split-sign-changes } ps \ x_0. \text{sign-changes } ps' \ x$
shows *eventually* $(\lambda x. \text{sign-changes}' \ x_0 \ ps \ x = \text{sign-changes } ps \ x)$ (at x_0)

<proof>

lemma (in *quasi-sturm-seq*) *hd-nonzero-imp-sign-changes-const-aux*:
assumes $\text{poly } (hd \ ps) \ x_0 \neq 0$ **and** $ps' \in \text{set } (\text{split-sign-changes } ps \ x_0)$
shows *eventually* $(\lambda x. \text{sign-changes } ps' \ x = \text{sign-changes } ps' \ x_0)$ (at x_0)

<proof>

lemma (in *quasi-sturm-seq*) *hd-nonzero-imp-sign-changes-const*:
assumes $\text{poly } (hd \ ps) \ x_0 \neq 0$
shows *eventually* $(\lambda x. \text{sign-changes } ps \ x = \text{sign-changes } ps \ x_0)$ (at x_0)

<proof>

lemma (in *sturm-seq*) *p-nonzero-imp-sign-changes-const*:
 $\text{poly } p \ x_0 \neq 0 \implies$
eventually $(\lambda x. \text{sign-changes } ps \ x = \text{sign-changes } ps \ x_0)$ (at x_0)

<proof>

If x is a root of p and p is not the zero polynomial, the number of sign changes of a Sturm chain of p decreases by 1 at x .

lemma (in *sturm-seq*) *p-zero*:
assumes $\text{poly } p \ x_0 = 0 \ p \neq 0$

shows *eventually* ($\lambda x. \text{sign-changes } ps \ x =$
 $\text{sign-changes } ps \ x_0 + (\text{if } x < x_0 \text{ then } 1 \text{ else } 0)$) (*at* x_0)
 ⟨*proof*⟩

With these two results, we can now show that if p is nonzero, the number of roots in an interval of the form $(a; b]$ is the difference of the sign changes of a Sturm sequence of p at a and b .

First, however, we prove the following auxiliary lemma that shows that if a function $f : \mathbb{R} \rightarrow \mathbb{N}$ is locally constant at any $x \in (a; b]$, it is constant across the entire interval $(a; b]$:

lemma *count-roots-between-aux*:

assumes $a \leq b$
assumes $\forall x::\text{real}. a < x \wedge x \leq b \longrightarrow \text{eventually } (\lambda \xi. f \ \xi = (f \ x::\text{nat}))$ (*at* x)
shows $\forall x. a < x \wedge x \leq b \longrightarrow f \ x = f \ b$
 ⟨*proof*⟩

Now we can prove the actual root-counting theorem:

theorem (*in sturm-seq*) *count-roots-between*:

assumes [*simp*]: $p \neq 0 \ a \leq b$
shows $\text{sign-changes } ps \ a - \text{sign-changes } ps \ b =$
 $\text{card } \{x. x > a \wedge x \leq b \wedge \text{poly } p \ x = 0\}$
 ⟨*proof*⟩

By applying this result to a sufficiently large upper bound, we can effectively count the number of roots “between a and infinity”, i. e. the roots greater than a :

lemma (*in sturm-seq*) *count-roots-above*:

assumes $p \neq 0$
shows $\text{sign-changes } ps \ a - \text{sign-changes-inf } ps =$
 $\text{card } \{x. x > a \wedge \text{poly } p \ x = 0\}$
 ⟨*proof*⟩

The same works analogously for the number of roots below a and the total number of roots.

lemma (*in sturm-seq*) *count-roots-below*:

assumes $p \neq 0$
shows $\text{sign-changes-neg-inf } ps - \text{sign-changes } ps \ a =$
 $\text{card } \{x. x \leq a \wedge \text{poly } p \ x = 0\}$
 ⟨*proof*⟩

lemma (*in sturm-seq*) *count-roots*:

assumes $p \neq 0$
shows $\text{sign-changes-neg-inf } ps - \text{sign-changes-inf } ps =$
 $\text{card } \{x. \text{poly } p \ x = 0\}$
 ⟨*proof*⟩

2.4 Constructing Sturm sequences

2.5 The canonical Sturm sequence

In this subsection, we will present the canonical Sturm sequence construction for a polynomial p without multiple roots that is very similar to the Euclidean algorithm:

$$p_i = \begin{cases} p & \text{for } i = 1 \\ p' & \text{for } i = 2 \\ -p_{i-2} \bmod p_{i-1} & \text{otherwise} \end{cases}$$

We break off the sequence at the first constant polynomial.

<proof>

function *sturm-aux* **where**

sturm-aux ($p :: \text{real poly}$) $q =$

(if degree $q = 0$ then $[p, q]$ else $p \# \text{sturm-aux } q \ (-(p \bmod q))$)

<proof>

termination *<proof>*

definition *sturm* **where** *sturm* $p = \text{sturm-aux } p \ (\text{pderiv } p)$

Next, we show some simple facts about this construction:

lemma *sturm-0[simp]*: *sturm* $0 = [0, 0]$

<proof>

lemma *[simp]*: *sturm-aux* $p \ q = [] \longleftrightarrow \text{False}$

<proof>

lemma *sturm-neq-Nil[simp]*: *sturm* $p \neq []$ *<proof>*

lemma *[simp]*: *hd* (*sturm* p) = p

<proof>

lemma *[simp]*: $p \in \text{set } (\text{sturm } p)$

<proof>

lemma *[simp]*: *length* (*sturm* p) ≥ 2

<proof>

lemma *[simp]*: *degree* (*last* (*sturm* p)) = 0

<proof>

lemma *[simp]*: *sturm-aux* $p \ q \ ! \ 0 = p$

<proof>

lemma *[simp]*: *sturm-aux* $p \ q \ ! \ \text{Suc } 0 = q$

<proof>

lemma *[simp]*: *sturm* $p \ ! \ 0 = p$

$\langle proof \rangle$
lemma [simp]: $sturm\ p\ !\ Suc\ 0 = pderiv\ p$
 $\langle proof \rangle$

lemma *sturm-indices*:
assumes $i < length\ (sturm\ p) - 2$
shows $sturm\ p!(i+2) = -(sturm\ p!i\ mod\ sturm\ p!(i+1))$
 $\langle proof \rangle$

If the Sturm sequence construction is applied to polynomials p and q , the greatest common divisor of p and q a divisor of every element in the sequence. This is obvious from the similarity to Euclid's algorithm for computing the GCD.

lemma *sturm-aux-gcd*: $r \in set\ (sturm\ aux\ p\ q) \implies gcd\ p\ q\ dvd\ r$
 $\langle proof \rangle$

lemma *sturm-gcd*: $r \in set\ (sturm\ p) \implies gcd\ p\ (pderiv\ p)\ dvd\ r$
 $\langle proof \rangle$

If two adjacent polynomials in the result of the canonical Sturm chain construction both have a root at some x , this x is a root of all polynomials in the sequence.

lemma *sturm-adjacent-root-propagate-left*:
assumes $i < length\ (sturm\ (p :: real\ poly)) - 1$
assumes $poly\ (sturm\ p\ !\ i)\ x = 0$
and $poly\ (sturm\ p\ !\ (i + 1))\ x = 0$
shows $\forall j \leq i+1. poly\ (sturm\ p\ !\ j)\ x = 0$
 $\langle proof \rangle$

Consequently, if this is the case in the canonical Sturm chain of p , p must have multiple roots.

lemma *sturm-adjacent-root-not-squarefree*:
assumes $i < length\ (sturm\ (p :: real\ poly)) - 1$
 $poly\ (sturm\ p\ !\ i)\ x = 0\ poly\ (sturm\ p\ !\ (i + 1))\ x = 0$
shows $\neg rsquarefree\ p$
 $\langle proof \rangle$

Since the second element of the sequence is chosen to be the derivative of p , p_1 and p_2 fulfil the property demanded by the definition of a Sturm sequence that they locally have opposite sign left of a root x of p and the same sign to the right of x .

lemma *sturm-firsttwo-signs-aux*:
assumes $(p :: real\ poly) \neq 0\ q \neq 0$
assumes $q\ pderiv$:
 $eventually\ (\lambda x. sgn\ (poly\ q\ x) = sgn\ (poly\ (pderiv\ p)\ x))\ (at\ x_0)$
assumes $p-0$: $poly\ p\ (x_0 :: real) = 0$

shows eventually $(\lambda x. \text{sgn} (\text{poly} (p * q) x) = (\text{if } x > x_0 \text{ then } 1 \text{ else } -1))$ (at x_0)
 <proof>

lemma *sturm-firsttwo-signs*:

fixes $ps :: \text{real poly list}$

assumes *squarefree*: $rsquarefree\ p$

assumes *p-0*: $\text{poly } p (x_0 :: \text{real}) = 0$

shows eventually $(\lambda x. \text{sgn} (\text{poly} (p * \text{sturm } p ! 1) x) =$
 $(\text{if } x > x_0 \text{ then } 1 \text{ else } -1))$ (at x_0)

<proof>

The construction also obviously fulfils the property about three adjacent polynomials in the sequence.

lemma *sturm-signs*:

assumes *squarefree*: $rsquarefree\ p$

assumes *i-in-range*: $i < \text{length} (\text{sturm } (p :: \text{real poly})) - 2$

assumes *q-0*: $\text{poly} (\text{sturm } p ! (i+1)) x = 0$ (**is** $\text{poly } ?q x = 0$)

shows $\text{poly} (\text{sturm } p ! (i+2)) x * \text{poly} (\text{sturm } p ! i) x < 0$
 (**is** $\text{poly } ?p x * \text{poly } ?r x < 0$)

<proof>

Finally, if p contains no multiple roots, *sturm* p , i.e. the canonical Sturm sequence for p , is a Sturm sequence and can be used to determine the number of roots of p .

lemma *sturm-seq-sturm[simp]*:

assumes *rsquarefree* p

shows *sturm-seq* $(\text{sturm } p) p$

<proof>

2.5.1 Canonical squarefree Sturm sequence

The previous construction does not work for polynomials with multiple roots, but we can simply “divide away” multiple roots by dividing p by the GCD of p and p' . The resulting polynomial has the same roots as p , but with multiplicity 1, allowing us to again use the canonical construction.

definition *sturm-squarefree* **where**

$\text{sturm-squarefree } p = \text{sturm } (p \text{ div } (\text{gcd } p (pderiv\ p)))$

lemma *sturm-squarefree-not-Nil[simp]*: $\text{sturm-squarefree } p \neq []$

<proof>

lemma *sturm-seq-sturm-squarefree*:

assumes [simp]: $p \neq 0$

defines [simp]: $p' \equiv p \text{ div } \text{gcd } p (pderiv\ p)$

shows *sturm-seq* $(\text{sturm-squarefree } p) p'$

<proof>

2.5.2 Optimisation for multiple roots

We can also define the following non-canonical Sturm sequence that is obtained by taking the canonical Sturm sequence of p (possibly with multiple roots) and then dividing the entire sequence by the GCD of p and its derivative.

definition *sturm-squarefree'* **where**

sturm-squarefree' $p = (\text{let } d = \text{gcd } p \text{ (pderiv } p)$
 in map $(\lambda p'. p' \text{ div } d) \text{ (sturm } p)$)

This construction also has all the desired properties:

lemma *sturm-squarefree'-adjacent-root-propagate-left*:

assumes $p \neq 0$
assumes $i < \text{length (sturm-squarefree' (p :: real poly))} - 1$
assumes $\text{poly (sturm-squarefree' p ! } i) x = 0$
and $\text{poly (sturm-squarefree' p ! (i + 1)) } x = 0$
shows $\forall j \leq i+1. \text{poly (sturm-squarefree' p ! } j) x = 0$
 $\langle \text{proof} \rangle$

lemma *sturm-squarefree'-adjacent-roots*:

assumes $p \neq 0$
 $i < \text{length (sturm-squarefree' (p :: real poly))} - 1$
 $\text{poly (sturm-squarefree' p ! } i) x = 0$
 $\text{poly (sturm-squarefree' p ! (i + 1)) } x = 0$
shows *False*
 $\langle \text{proof} \rangle$

lemma *sturm-squarefree'-signs*:

assumes $p \neq 0$
assumes *i-in-range*: $i < \text{length (sturm-squarefree' (p :: real poly))} - 2$
assumes *q-0*: $\text{poly (sturm-squarefree' p ! (i+1)) } x = 0$ (**is** $\text{poly } ?q x = 0$)
shows $\text{poly (sturm-squarefree' p ! (i+2)) } x *$
 $\text{poly (sturm-squarefree' p ! } i) x < 0$
 (**is** $\text{poly } ?r x * \text{poly } ?p x < 0$)
 $\langle \text{proof} \rangle$

This approach indeed also yields a valid squarefree Sturm sequence for the polynomial $p/\text{gcd}(p, p')$.

lemma *sturm-seq-sturm-squarefree'*:

assumes $(p :: \text{real poly}) \neq 0$
defines $d \equiv \text{gcd } p \text{ (pderiv } p)$
shows $\text{sturm-seq (sturm-squarefree' } p) (p \text{ div } d)$
 (**is** $\text{sturm-seq } ?ps' ?p'$)
 $\langle \text{proof} \rangle$

This construction is obviously more expensive to compute than the one that *first* divides p by $\text{gcd}(p, p')$ and *then* applies the canonical construction. In this construction, we *first* compute the canonical Sturm sequence of p as

if it had no multiple roots and *then* divide by the GCD. However, it can be seen quite easily that unless x is a multiple root of p , i.e. as long as $\gcd(P, P') \neq 0$, the number of sign changes in a sequence of polynomials does not actually change when we divide the polynomials by $\gcd(p, p')$. Therefore we can use the canonical Sturm sequence even in the non-square-free case as long as the borders of the interval we are interested in are not multiple roots of the polynomial.

lemma *sign-changes-mult-aux*:

assumes $d \neq (0::real)$

shows $length\ (remdups-adj\ (filter\ (\lambda x. x \neq 0)\ (map\ ((*)\ d \circ f)\ xs))) =$
 $length\ (remdups-adj\ (filter\ (\lambda x. x \neq 0)\ (map\ f\ xs)))$

<proof>

lemma *sturm-sturm-squarefree'-same-sign-changes*:

fixes $p :: real\ poly$

defines $ps \equiv sturm\ p$ **and** $ps' \equiv sturm-squarefree'\ p$

shows $poly\ p\ x \neq 0 \vee poly\ (pderiv\ p)\ x \neq 0 \implies$

$sign-changes\ ps'\ x = sign-changes\ ps\ x$

$p \neq 0 \implies sign-changes-inf\ ps' = sign-changes-inf\ ps$

$p \neq 0 \implies sign-changes-neg-inf\ ps' = sign-changes-neg-inf\ ps$

<proof>

2.6 Root-counting functions

With all these results, we can now define functions that count roots in bounded and unbounded intervals:

definition *count-roots-between* **where**

count-roots-between $p\ a\ b = (if\ a \leq b \wedge p \neq 0\ then$

$(let\ ps = sturm-squarefree\ p$

$in\ sign-changes\ ps\ a - sign-changes\ ps\ b)$ *else* $0)$

definition *count-roots* **where**

count-roots $p = (if\ (p::real\ poly) = 0\ then\ 0\ else$

$(let\ ps = sturm-squarefree\ p$

$in\ sign-changes-neg-inf\ ps - sign-changes-inf\ ps))$

definition *count-roots-above* **where**

count-roots-above $p\ a = (if\ (p::real\ poly) = 0\ then\ 0\ else$

$(let\ ps = sturm-squarefree\ p$

$in\ sign-changes\ ps\ a - sign-changes-inf\ ps))$

definition *count-roots-below* **where**

count-roots-below $p\ a = (if\ (p::real\ poly) = 0\ then\ 0\ else$

$(let\ ps = sturm-squarefree\ p$

$in\ sign-changes-neg-inf\ ps - sign-changes\ ps\ a))$

lemma *count-roots-between-correct*:

count-roots-between p a $b = \text{card } \{x. a < x \wedge x \leq b \wedge \text{poly } p x = 0\}$
<proof>

lemma *count-roots-correct*:

fixes $p :: \text{real poly}$
shows *count-roots* $p = \text{card } \{x. \text{poly } p x = 0\}$ (**is** $- = \text{card } ?S$)
<proof>

lemma *count-roots-above-correct*:

fixes $p :: \text{real poly}$
shows *count-roots-above* p $a = \text{card } \{x. x > a \wedge \text{poly } p x = 0\}$
(**is** $- = \text{card } ?S$)
<proof>

lemma *count-roots-below-correct*:

fixes $p :: \text{real poly}$
shows *count-roots-below* p $a = \text{card } \{x. x \leq a \wedge \text{poly } p x = 0\}$
(**is** $- = \text{card } ?S$)
<proof>

The optimisation explained above can be used to prove more efficient code equations that use the more efficient construction in the case that the interval borders are not multiple roots:

lemma *count-roots-between[code]*:

count-roots-between p a $b =$
 (*let* $q = \text{pderiv } p$
 in if $a > b \vee p = 0$ *then* 0
 else if $(\text{poly } p a \neq 0 \vee \text{poly } q a \neq 0) \wedge (\text{poly } p b \neq 0 \vee \text{poly } q b \neq 0)$
 then (*let* $ps = \text{sturm } p$
 in $\text{sign-changes } ps a - \text{sign-changes } ps b$)
 else (*let* $ps = \text{sturm-squarefree } p$
 in $\text{sign-changes } ps a - \text{sign-changes } ps b$)
<proof>

lemma *count-roots-code[code]*:

count-roots $(p::\text{real poly}) =$
 (*if* $p = 0$ *then* 0
 else let $ps = \text{sturm } p$
 in $\text{sign-changes-neg-inf } ps - \text{sign-changes-inf } ps$)
<proof>

lemma *count-roots-above-code[code]*:

count-roots-above p $a =$
 (*let* $q = \text{pderiv } p$
 in if $p = 0$ *then* 0
 else if $\text{poly } p a \neq 0 \vee \text{poly } q a \neq 0$

```

    then (let ps = sturm p
          in sign-changes ps a - sign-changes-inf ps)
    else (let ps = sturm-squarefree p
          in sign-changes ps a - sign-changes-inf ps))
⟨proof⟩

```

lemma *count-roots-below-code*[code]:

```

count-roots-below p a =
  (let q = pderiv p
   in if p = 0 then 0
   else if poly p a ≠ 0 ∨ poly q a ≠ 0
        then (let ps = sturm p
              in sign-changes-neg-inf ps - sign-changes ps a)
        else (let ps = sturm-squarefree p
              in sign-changes-neg-inf ps - sign-changes ps a))
⟨proof⟩

```

end

3 The “sturm” proof method

```

theory Sturm-Method
imports Sturm-Theorem
begin

```

3.1 Preliminary lemmas

In this subsection, we prove lemmas that reduce root counting and related statements to simple, computable expressions using the *count-roots* function family.

lemma *poly-card-roots-less-leq*:

```

card {x. a < x ∧ x ≤ b ∧ poly p x = 0} = count-roots-between p a b
⟨proof⟩

```

lemma *poly-card-roots-leq-leq*:

```

card {x. a ≤ x ∧ x ≤ b ∧ poly p x = 0} =
  ( count-roots-between p a b +
    (if (a ≤ b ∧ poly p a = 0 ∧ p ≠ 0) ∨ (a = b ∧ p = 0) then 1 else 0))
⟨proof⟩

```

lemma *poly-card-roots-less-less*:

```

card {x. a < x ∧ x < b ∧ poly p x = 0} =
  ( count-roots-between p a b -
    (if poly p b = 0 ∧ a < b ∧ p ≠ 0 then 1 else 0))
⟨proof⟩

```

lemma *poly-card-roots-leq-less*:

```

card {x::real. a ≤ x ∧ x < b ∧ poly p x = 0} =

```

$(\text{count-roots-between } p \ a \ b +$
 $(\text{if } p \neq 0 \wedge a < b \wedge \text{poly } p \ a = 0 \text{ then } 1 \text{ else } 0) -$
 $(\text{if } p \neq 0 \wedge a < b \wedge \text{poly } p \ b = 0 \text{ then } 1 \text{ else } 0))$
 <proof>

lemma *poly-card-roots*:
 $\text{card } \{x::\text{real}. \text{poly } p \ x = 0\} = \text{count-roots } p$
 <proof>

lemma *poly-no-roots*:
 $(\forall x. \text{poly } p \ x \neq 0) \iff (p \neq 0 \wedge \text{count-roots } p = 0)$
 <proof>

lemma *poly-pos*:
 $(\forall x. \text{poly } p \ x > 0) \iff ($
 $p \neq 0 \wedge \text{poly-inf } p = 1 \wedge \text{count-roots } p = 0)$
 <proof>

lemma *poly-card-roots-greater*:
 $\text{card } \{x::\text{real}. x > a \wedge \text{poly } p \ x = 0\} = \text{count-roots-above } p \ a$
 <proof>

lemma *poly-card-roots-leq*:
 $\text{card } \{x::\text{real}. x \leq a \wedge \text{poly } p \ x = 0\} = \text{count-roots-below } p \ a$
 <proof>

lemma *poly-card-roots-geq*:
 $\text{card } \{x::\text{real}. x \geq a \wedge \text{poly } p \ x = 0\} = ($
 $\text{count-roots-above } p \ a + (\text{if } \text{poly } p \ a = 0 \wedge p \neq 0 \text{ then } 1 \text{ else } 0))$
 <proof>

lemma *poly-card-roots-less*:
 $\text{card } \{x::\text{real}. x < a \wedge \text{poly } p \ x = 0\} =$
 $(\text{count-roots-below } p \ a - (\text{if } \text{poly } p \ a = 0 \wedge p \neq 0 \text{ then } 1 \text{ else } 0))$
 <proof>

lemma *poly-no-roots-less-leq*:
 $(\forall x. a < x \wedge x \leq b \implies \text{poly } p \ x \neq 0) \iff$
 $((a \geq b \vee (p \neq 0 \wedge \text{count-roots-between } p \ a \ b = 0)))$
 <proof>

lemma *poly-pos-between-less-leq*:
 $(\forall x. a < x \wedge x \leq b \implies \text{poly } p \ x > 0) \iff$
 $((a \geq b \vee (p \neq 0 \wedge \text{poly } p \ b > 0 \wedge \text{count-roots-between } p \ a \ b = 0)))$
 <proof>

lemma *poly-no-roots-leq-leq*:

$$(\forall x. a \leq x \wedge x \leq b \longrightarrow \text{poly } p \ x \neq 0) \longleftrightarrow$$

$$((a > b \vee (p \neq 0 \wedge \text{poly } p \ a \neq 0 \wedge \text{count-roots-between } p \ a \ b = 0)))$$

<proof>

lemma *poly-pos-between-leq-leq*:

$$(\forall x. a \leq x \wedge x \leq b \longrightarrow \text{poly } p \ x > 0) \longleftrightarrow$$

$$((a > b \vee (p \neq 0 \wedge \text{poly } p \ a > 0 \wedge$$

$$\text{count-roots-between } p \ a \ b = 0)))$$

<proof>

lemma *poly-no-roots-less-less*:

$$(\forall x. a < x \wedge x < b \longrightarrow \text{poly } p \ x \neq 0) \longleftrightarrow$$

$$((a \geq b \vee p \neq 0 \wedge \text{count-roots-between } p \ a \ b =$$

$$(\text{if } \text{poly } p \ b = 0 \text{ then } 1 \text{ else } 0)))$$

<proof>

lemma *poly-pos-between-less-less*:

$$(\forall x. a < x \wedge x < b \longrightarrow \text{poly } p \ x > 0) \longleftrightarrow$$

$$((a \geq b \vee (p \neq 0 \wedge \text{poly } p \ ((a+b)/2) > 0 \wedge$$

$$\text{count-roots-between } p \ a \ b = (\text{if } \text{poly } p \ b = 0 \text{ then } 1 \text{ else } 0))))$$

<proof>

lemma *poly-no-roots-leq-less*:

$$(\forall x. a \leq x \wedge x < b \longrightarrow \text{poly } p \ x \neq 0) \longleftrightarrow$$

$$((a \geq b \vee p \neq 0 \wedge \text{poly } p \ a \neq 0 \wedge \text{count-roots-between } p \ a \ b =$$

$$(\text{if } a < b \wedge \text{poly } p \ b = 0 \text{ then } 1 \text{ else } 0)))$$

<proof>

lemma *poly-pos-between-leq-less*:

$$(\forall x. a \leq x \wedge x < b \longrightarrow \text{poly } p \ x > 0) \longleftrightarrow$$

$$((a \geq b \vee (p \neq 0 \wedge \text{poly } p \ a > 0 \wedge \text{count-roots-between } p \ a \ b =$$

$$(\text{if } a < b \wedge \text{poly } p \ b = 0 \text{ then } 1 \text{ else } 0))))$$

<proof>

lemma *poly-no-roots-greater*:

$$(\forall x. x > a \longrightarrow \text{poly } p \ x \neq 0) \longleftrightarrow$$

$$((p \neq 0 \wedge \text{count-roots-above } p \ a = 0))$$

<proof>

lemma *poly-pos-greater*:

$$(\forall x. x > a \longrightarrow \text{poly } p \ x > 0) \longleftrightarrow ($$

$$p \neq 0 \wedge \text{poly-inf } p = 1 \wedge \text{count-roots-above } p \ a = 0)$$

<proof>

lemma *poly-no-roots-leq*:

$(\forall x. x \leq a \longrightarrow \text{poly } p \ x \neq 0) \longleftrightarrow$
 $(p \neq 0 \wedge \text{count-roots-below } p \ a = 0)$
 <proof>

lemma *poly-pos-leq*:

$(\forall x. x \leq a \longrightarrow \text{poly } p \ x > 0) \longleftrightarrow$
 $(p \neq 0 \wedge \text{poly-neg-inf } p = 1 \wedge \text{count-roots-below } p \ a = 0)$
 <proof>

lemma *poly-no-roots-geq*:

$(\forall x. x \geq a \longrightarrow \text{poly } p \ x \neq 0) \longleftrightarrow$
 $(p \neq 0 \wedge \text{poly } p \ a \neq 0 \wedge \text{count-roots-above } p \ a = 0)$
 <proof>

lemma *poly-pos-geq*:

$(\forall x. x \geq a \longrightarrow \text{poly } p \ x > 0) \longleftrightarrow$
 $(p \neq 0 \wedge \text{poly-inf } p = 1 \wedge \text{poly } p \ a \neq 0 \wedge \text{count-roots-above } p \ a = 0)$
 <proof>

lemma *poly-no-roots-less*:

$(\forall x. x < a \longrightarrow \text{poly } p \ x \neq 0) \longleftrightarrow$
 $((p \neq 0 \wedge \text{count-roots-below } p \ a = (\text{if } \text{poly } p \ a = 0 \text{ then } 1 \text{ else } 0)))$
 <proof>

lemma *poly-pos-less*:

$(\forall x. x < a \longrightarrow \text{poly } p \ x > 0) \longleftrightarrow$
 $(p \neq 0 \wedge \text{poly-neg-inf } p = 1 \wedge \text{count-roots-below } p \ a =$
 $(\text{if } \text{poly } p \ a = 0 \text{ then } 1 \text{ else } 0))$
 <proof>

lemmas *sturm-card-substs* = *poly-card-roots poly-card-roots-less-leq*
poly-card-roots-leq-less poly-card-roots-less-less poly-card-roots-leq-leq
poly-card-roots-less poly-card-roots-leq poly-card-roots-greater
poly-card-roots-geq

lemmas *sturm-prop-substs* = *poly-no-roots poly-no-roots-less-leq*
poly-no-roots-leq-leq poly-no-roots-less-less poly-no-roots-leq-less
poly-no-roots-leq poly-no-roots-less poly-no-roots-geq
poly-no-roots-greater
poly-pos poly-pos-greater poly-pos-geq poly-pos-less poly-pos-leq
poly-pos-between-leq-less poly-pos-between-less-leq
poly-pos-between-leq-leq poly-pos-between-less-less

3.2 Reification

This subsection defines a number of equations to automatically convert statements about roots of polynomials into a canonical form so that they can be proven using the above substitutions.

definition $PR\text{-}TAG\ x \equiv x$

lemma $sturm\text{-}id\text{-}PR\text{-}prio0$:

$$\begin{aligned} \{x::real. P\ x\} &= \{x::real. (PR\text{-}TAG\ P)\ x\} \\ (\forall x::real. f\ x < g\ x) &= (\forall x::real. PR\text{-}TAG\ (\lambda x. f\ x < g\ x)\ x) \\ (\forall x::real. P\ x) &= (\forall x::real. \neg(PR\text{-}TAG\ (\lambda x. \neg P\ x))\ x) \\ &\langle proof \rangle \end{aligned}$$

lemma $sturm\text{-}id\text{-}PR\text{-}prio1$:

$$\begin{aligned} \{x::real. x < a \wedge P\ x\} &= \{x::real. x < a \wedge (PR\text{-}TAG\ P)\ x\} \\ \{x::real. x \leq a \wedge P\ x\} &= \{x::real. x \leq a \wedge (PR\text{-}TAG\ P)\ x\} \\ \{x::real. x \geq b \wedge P\ x\} &= \{x::real. x \geq b \wedge (PR\text{-}TAG\ P)\ x\} \\ \{x::real. x > b \wedge P\ x\} &= \{x::real. x > b \wedge (PR\text{-}TAG\ P)\ x\} \\ (\forall x::real < a. f\ x < g\ x) &= (\forall x::real < a. PR\text{-}TAG\ (\lambda x. f\ x < g\ x)\ x) \\ (\forall x::real \leq a. f\ x < g\ x) &= (\forall x::real \leq a. PR\text{-}TAG\ (\lambda x. f\ x < g\ x)\ x) \\ (\forall x::real > a. f\ x < g\ x) &= (\forall x::real > a. PR\text{-}TAG\ (\lambda x. f\ x < g\ x)\ x) \\ (\forall x::real \geq a. f\ x < g\ x) &= (\forall x::real \geq a. PR\text{-}TAG\ (\lambda x. f\ x < g\ x)\ x) \\ (\forall x::real < a. P\ x) &= (\forall x::real < a. \neg(PR\text{-}TAG\ (\lambda x. \neg P\ x))\ x) \\ (\forall x::real > a. P\ x) &= (\forall x::real > a. \neg(PR\text{-}TAG\ (\lambda x. \neg P\ x))\ x) \\ (\forall x::real \leq a. P\ x) &= (\forall x::real \leq a. \neg(PR\text{-}TAG\ (\lambda x. \neg P\ x))\ x) \\ (\forall x::real \geq a. P\ x) &= (\forall x::real \geq a. \neg(PR\text{-}TAG\ (\lambda x. \neg P\ x))\ x) \\ &\langle proof \rangle \end{aligned}$$

lemma $sturm\text{-}id\text{-}PR\text{-}prio2$:

$$\begin{aligned} \{x::real. x > a \wedge x \leq b \wedge P\ x\} &= \\ &\{x::real. x > a \wedge x \leq b \wedge PR\text{-}TAG\ P\ x\} \\ \{x::real. x \geq a \wedge x \leq b \wedge P\ x\} &= \\ &\{x::real. x \geq a \wedge x \leq b \wedge PR\text{-}TAG\ P\ x\} \\ \{x::real. x \geq a \wedge x < b \wedge P\ x\} &= \\ &\{x::real. x \geq a \wedge x < b \wedge PR\text{-}TAG\ P\ x\} \\ \{x::real. x > a \wedge x < b \wedge P\ x\} &= \\ &\{x::real. x > a \wedge x < b \wedge PR\text{-}TAG\ P\ x\} \\ (\forall x::real. a < x \wedge x \leq b \longrightarrow f\ x < g\ x) &= \\ &(\forall x::real. a < x \wedge x \leq b \longrightarrow PR\text{-}TAG\ (\lambda x. f\ x < g\ x)\ x) \\ (\forall x::real. a \leq x \wedge x \leq b \longrightarrow f\ x < g\ x) &= \\ &(\forall x::real. a \leq x \wedge x \leq b \longrightarrow PR\text{-}TAG\ (\lambda x. f\ x < g\ x)\ x) \\ (\forall x::real. a < x \wedge x < b \longrightarrow f\ x < g\ x) &= \\ &(\forall x::real. a < x \wedge x < b \longrightarrow PR\text{-}TAG\ (\lambda x. f\ x < g\ x)\ x) \\ (\forall x::real. a \leq x \wedge x < b \longrightarrow f\ x < g\ x) &= \\ &(\forall x::real. a \leq x \wedge x < b \longrightarrow PR\text{-}TAG\ (\lambda x. f\ x < g\ x)\ x) \\ (\forall x::real. a < x \wedge x \leq b \longrightarrow P\ x) &= \\ &(\forall x::real. a < x \wedge x \leq b \longrightarrow \neg(PR\text{-}TAG\ (\lambda x. \neg P\ x))\ x) \\ (\forall x::real. a \leq x \wedge x \leq b \longrightarrow P\ x) &= \\ &(\forall x::real. a \leq x \wedge x \leq b \longrightarrow \neg(PR\text{-}TAG\ (\lambda x. \neg P\ x))\ x) \end{aligned}$$

$(\forall x::real. a \leq x \wedge x < b \longrightarrow P x) =$
 $(\forall x::real. a \leq x \wedge x < b \longrightarrow \neg(PR-TAG (\lambda x. \neg P x)) x)$
 $(\forall x::real. a < x \wedge x < b \longrightarrow P x) =$
 $(\forall x::real. a < x \wedge x < b \longrightarrow \neg(PR-TAG (\lambda x. \neg P x)) x)$
 $\langle proof \rangle$

lemma *PR-TAG-intro-prio0*:

fixes $P :: real \Rightarrow bool$ **and** $f :: real \Rightarrow real$

shows

$PR-TAG P = P' \Longrightarrow PR-TAG (\lambda x. \neg(\neg P x)) = P'$
 $\llbracket PR-TAG P = (\lambda x. poly p x = 0); PR-TAG Q = (\lambda x. poly q x = 0) \rrbracket$
 $\Longrightarrow PR-TAG (\lambda x. P x \wedge Q x) = (\lambda x. poly (gcd p q) x = 0)$ **and**
 $\llbracket PR-TAG P = (\lambda x. poly p x = 0); PR-TAG Q = (\lambda x. poly q x = 0) \rrbracket$
 $\Longrightarrow PR-TAG (\lambda x. P x \vee Q x) = (\lambda x. poly (p*q) x = 0)$ **and**

$\llbracket PR-TAG f = (\lambda x. poly p x); PR-TAG g = (\lambda x. poly q x) \rrbracket$
 $\Longrightarrow PR-TAG (\lambda x. f x = g x) = (\lambda x. poly (p-q) x = 0)$
 $\llbracket PR-TAG f = (\lambda x. poly p x); PR-TAG g = (\lambda x. poly q x) \rrbracket$
 $\Longrightarrow PR-TAG (\lambda x. f x \neq g x) = (\lambda x. poly (p-q) x \neq 0)$
 $\llbracket PR-TAG f = (\lambda x. poly p x); PR-TAG g = (\lambda x. poly q x) \rrbracket$
 $\Longrightarrow PR-TAG (\lambda x. f x < g x) = (\lambda x. poly (q-p) x > 0)$
 $\llbracket PR-TAG f = (\lambda x. poly p x); PR-TAG g = (\lambda x. poly q x) \rrbracket$
 $\Longrightarrow PR-TAG (\lambda x. f x \leq g x) = (\lambda x. poly (q-p) x \geq 0)$

$PR-TAG f = (\lambda x. poly p x) \Longrightarrow PR-TAG (\lambda x. -f x) = (\lambda x. poly (-p) x)$
 $\llbracket PR-TAG f = (\lambda x. poly p x); PR-TAG g = (\lambda x. poly q x) \rrbracket$
 $\Longrightarrow PR-TAG (\lambda x. f x + g x) = (\lambda x. poly (p+q) x)$
 $\llbracket PR-TAG f = (\lambda x. poly p x); PR-TAG g = (\lambda x. poly q x) \rrbracket$
 $\Longrightarrow PR-TAG (\lambda x. f x - g x) = (\lambda x. poly (p-q) x)$
 $\llbracket PR-TAG f = (\lambda x. poly p x); PR-TAG g = (\lambda x. poly q x) \rrbracket$
 $\Longrightarrow PR-TAG (\lambda x. f x * g x) = (\lambda x. poly (p*q) x)$
 $PR-TAG f = (\lambda x. poly p x) \Longrightarrow PR-TAG (\lambda x. (f x) \widehat{n}) = (\lambda x. poly (p \widehat{n}) x)$
 $PR-TAG (\lambda x. poly p x :: real) = (\lambda x. poly p x)$
 $PR-TAG (\lambda x. x::real) = (\lambda x. poly [0,1:] x)$
 $PR-TAG (\lambda x. a::real) = (\lambda x. poly [a:] x)$
 $\langle proof \rangle$

lemma *PR-TAG-intro-prio1*:

fixes $f :: real \Rightarrow real$

shows

$PR-TAG f = (\lambda x. poly p x) \Longrightarrow PR-TAG (\lambda x. f x = 0) = (\lambda x. poly p x = 0)$
 $PR-TAG f = (\lambda x. poly p x) \Longrightarrow PR-TAG (\lambda x. f x \neq 0) = (\lambda x. poly p x \neq 0)$
 $PR-TAG f = (\lambda x. poly p x) \Longrightarrow PR-TAG (\lambda x. 0 = f x) = (\lambda x. poly p x = 0)$
 $PR-TAG f = (\lambda x. poly p x) \Longrightarrow PR-TAG (\lambda x. 0 \neq f x) = (\lambda x. poly p x \neq 0)$
 $PR-TAG f = (\lambda x. poly p x) \Longrightarrow PR-TAG (\lambda x. f x \geq 0) = (\lambda x. poly p x \geq 0)$
 $PR-TAG f = (\lambda x. poly p x) \Longrightarrow PR-TAG (\lambda x. f x > 0) = (\lambda x. poly p x > 0)$

$PR\text{-TAG } f = (\lambda x. \text{poly } p \ x) \implies PR\text{-TAG } (\lambda x. f \ x \leq 0) = (\lambda x. \text{poly } (-p) \ x \geq 0)$
 $PR\text{-TAG } f = (\lambda x. \text{poly } p \ x) \implies PR\text{-TAG } (\lambda x. f \ x < 0) = (\lambda x. \text{poly } (-p) \ x > 0)$
 $PR\text{-TAG } f = (\lambda x. \text{poly } p \ x) \implies$
 $\quad PR\text{-TAG } (\lambda x. 0 \leq f \ x) = (\lambda x. \text{poly } (-p) \ x \leq 0)$
 $PR\text{-TAG } f = (\lambda x. \text{poly } p \ x) \implies$
 $\quad PR\text{-TAG } (\lambda x. 0 < f \ x) = (\lambda x. \text{poly } (-p) \ x < 0)$
 $PR\text{-TAG } f = (\lambda x. \text{poly } p \ x)$
 $\implies PR\text{-TAG } (\lambda x. a * f \ x) = (\lambda x. \text{poly } (\text{smult } a \ p) \ x)$
 $PR\text{-TAG } f = (\lambda x. \text{poly } p \ x)$
 $\implies PR\text{-TAG } (\lambda x. f \ x * a) = (\lambda x. \text{poly } (\text{smult } a \ p) \ x)$
 $PR\text{-TAG } f = (\lambda x. \text{poly } p \ x)$
 $\implies PR\text{-TAG } (\lambda x. f \ x / a) = (\lambda x. \text{poly } (\text{smult } (\text{inverse } a) \ p) \ x)$
 $PR\text{-TAG } (\lambda x. x \hat{n} :: \text{real}) = (\lambda x. \text{poly } (\text{monom } 1 \ n) \ x)$
 $\langle \text{proof} \rangle$

lemma *PR-TAG-intro-prio2*:

$PR\text{-TAG } (\lambda x. 1 / b) = (\lambda x. \text{inverse } b)$
 $PR\text{-TAG } (\lambda x. a / b) = (\lambda x. a / b)$
 $PR\text{-TAG } (\lambda x. a / b * x \hat{n} :: \text{real}) = (\lambda x. \text{poly } (\text{monom } (a/b) \ n) \ x)$
 $PR\text{-TAG } (\lambda x. x \hat{n} * a / b :: \text{real}) = (\lambda x. \text{poly } (\text{monom } (a/b) \ n) \ x)$
 $PR\text{-TAG } (\lambda x. a * x \hat{n} :: \text{real}) = (\lambda x. \text{poly } (\text{monom } a \ n) \ x)$
 $PR\text{-TAG } (\lambda x. x \hat{n} * a :: \text{real}) = (\lambda x. \text{poly } (\text{monom } a \ n) \ x)$
 $PR\text{-TAG } (\lambda x. x \hat{n} / a :: \text{real}) = (\lambda x. \text{poly } (\text{monom } (\text{inverse } a) \ n) \ x)$

$PR\text{-TAG } (\lambda x. f \ x \hat{ } (\text{Suc } (\text{Suc } 0)) :: \text{real}) = (\lambda x. \text{poly } p \ x)$
 $\implies PR\text{-TAG } (\lambda x. f \ x * f \ x :: \text{real}) = (\lambda x. \text{poly } p \ x)$
 $PR\text{-TAG } (\lambda x. (f \ x) \hat{ } \text{Suc } n :: \text{real}) = (\lambda x. \text{poly } p \ x)$
 $\implies PR\text{-TAG } (\lambda x. (f \ x) \hat{n} * f \ x :: \text{real}) = (\lambda x. \text{poly } p \ x)$
 $PR\text{-TAG } (\lambda x. (f \ x) \hat{ } \text{Suc } n :: \text{real}) = (\lambda x. \text{poly } p \ x)$
 $\implies PR\text{-TAG } (\lambda x. f \ x * (f \ x) \hat{n} :: \text{real}) = (\lambda x. \text{poly } p \ x)$
 $PR\text{-TAG } (\lambda x. (f \ x) \hat{(m+n)} :: \text{real}) = (\lambda x. \text{poly } p \ x)$
 $\implies PR\text{-TAG } (\lambda x. (f \ x) \hat{m} * (f \ x) \hat{n} :: \text{real}) = (\lambda x. \text{poly } p \ x)$
 $\langle \text{proof} \rangle$

lemma *sturm-meta-spec*: $(\bigwedge x :: \text{real}. P \ x) \implies P \ x$ $\langle \text{proof} \rangle$

lemma *sturm-imp-conv*:

$(a < x \longrightarrow x < b \longrightarrow c) \longleftrightarrow (a < x \wedge x < b \longrightarrow c)$
 $(a \leq x \longrightarrow x < b \longrightarrow c) \longleftrightarrow (a \leq x \wedge x < b \longrightarrow c)$
 $(a < x \longrightarrow x \leq b \longrightarrow c) \longleftrightarrow (a < x \wedge x \leq b \longrightarrow c)$
 $(a \leq x \longrightarrow x \leq b \longrightarrow c) \longleftrightarrow (a \leq x \wedge x \leq b \longrightarrow c)$
 $(x < b \longrightarrow a < x \longrightarrow c) \longleftrightarrow (a < x \wedge x < b \longrightarrow c)$
 $(x < b \longrightarrow a \leq x \longrightarrow c) \longleftrightarrow (a \leq x \wedge x < b \longrightarrow c)$
 $(x \leq b \longrightarrow a < x \longrightarrow c) \longleftrightarrow (a < x \wedge x \leq b \longrightarrow c)$
 $(x \leq b \longrightarrow a \leq x \longrightarrow c) \longleftrightarrow (a \leq x \wedge x \leq b \longrightarrow c)$
 $\langle \text{proof} \rangle$

3.3 Setup for the “sturm” method

$\langle ML \rangle$

end

```
theory Sturm
imports Sturm-Method
begin
```

end

4 Example usage of the “sturm” method

```
theory Sturm-Ex
imports ../Sturm
begin
```

In this section, we give a variety of statements about real polynomials that can be proven by the *sturm* method.

```
lemma
   $\forall x::real. x^2 + 1 \neq 0$ 
  <proof>
```

```
lemma
  fixes  $x :: real$ 
  shows  $x^2 + 1 \neq 0$  <proof>
```

```
lemma  $(x::real) > 1 \implies x^3 > 1$  <proof>
```

```
lemma  $\forall x::real. x*x \neq -1$  <proof>
```

schematic-goal A:

```
card { $x::real. -0.010831 < x \wedge x < 0.010831 \wedge$ 
   $1/120*x^5 + 1/24*x^4 + 1/6*x^3 - 49/16777216*x^2 - 17/2097152*x =$ 
  0}
  = ?n
  <proof>
```

```
lemma card { $x::real. x^3 + x = 2*x^2 \wedge x^3 - 6*x^2 + 11*x = 6$ } = 1
  <proof>
```

```
schematic-goal card { $x::real. x^3 + x = 2*x^2 \vee x^3 - 6*x^2 + 11*x = 6$ }
  = ?n <proof>
```

```
lemma
  card { $x::real. -0.010831 < x \wedge x < 0.010831 \wedge$ 
  poly [0, -17/2097152, -49/16777216, 1/6, 1/24, 1/120:]  $x = 0$ } = 3
  <proof>
```

```
lemma  $\forall x::real. x*x \neq 0 \vee x*x - 1 \neq 2*x$  <proof>
```

lemma $(x::real)*x+1 \neq 0 \wedge (x^2+1)*(x^2+2) \neq 0$ *(proof)*

3 examples related to continued fraction approximants to exp: LCP

lemma fixes $x::real$

shows $-7.29347719 \leq x \implies 0 < x^5 + 30*x^4 + 420*x^3 + 3360*x^2 + 15120*x + 30240$

(proof)

lemma fixes $x::real$

shows $0 < x^6 + 42*x^5 + 840*x^4 + 10080*x^3 + 75600*x^2 + 332640*x + 665280$

(proof)

schematic-goal card $\{x::real. x^7 + 56*x^6 + 1512*x^5 + 25200*x^4 + 277200*x^3 + 1995840*x^2 + 8648640*x = -17297280\} = ?n$

(proof)

end