Safe OCL

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Abstract

The theory is a formalization of the OCL type system, its abstract syntax and expression typing rules [1]. The theory does not define a concrete syntax and a semantics. In contrast to Featherweight OCL [2], it is based on a deep embedding approach. The type system is defined from scratch, it is not based on the Isabelle HOL type system.

The Safe OCL distincts nullable and non-nullable types. Also the theory gives a formal definition of safe navigation operations [3]. The Safe OCL typing rules are much stricter than rules given in the OCL specification. It allows one to catch more errors on a type checking phase.

The type theory presented is four-layered: classes, basic types, generic types, errorable types. We introduce the following new types: non-nullable types ($\tau[1]$), nullable types ($\tau[?]$), *OclSuper*. *OclSuper* is a supertype of all other types (basic types, collections, tuples). This type allows us to define a total supremum function, so types form an upper semilattice. It allows us to define rich expression typing rules in an elegant manner.

The Preliminaries Section of the theory defines a number of helper lemmas for transitive closures and tuples. It defines also a generic object model independent from OCL. It allows one to use the theory as a reference for formalization of analogous languages.

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Chapter 1

Preliminaries

1.1 Errorable

theory *Errorable*
imports *Main*
begin

notation *bot* (‹⊥›)

typedef *'a errorable* (‹-⊥› [*21*] *21*) = *UNIV* :: *'a option set* **..**

definition *errorable* :: *'a* ⇒ *'a errorable* (‹-⊥› [*1000*] *1000*) **where**
errorable x = *Abs-errorable* (*Some x*)

instantiation *errorable* :: (*type*) *bot*
begin
definition ⊥ ≡ *Abs-errorable None*
instance ..
end

free-constructors *case-errorable* **for**
errorable
| ⊥ :: *'a errorable*
unfolding *errorable-def bot-errorable-def*
apply (*metis Abs-errorable-cases not-None-eq*)
apply (*metis Abs-errorable-inverse UNIV-I option.inject*)
by (*simp add*: *Abs-errorable-inject*)

copy-bnf *'a errorable*

end

1.2 Transitive Closures

theory *Transitive-Closure-Ext*

imports *HOL−Library.FuncSet*
begin

1.2.1 Basic Properties

R^{++} is a transitive closure of a relation R. R^{**} is a reflexive transitive closure of a relation R.

A function f is surjective on R^{++} iff for any two elements in the range of f, related through R^{++}, all their intermediate elements belong to the range of f.

abbreviation *surj-on-trancl R f* ≡
(∀ *x y z*. R^{++} (*f x*) *y* ⟶ *R y* (*f z*) ⟶ *y* ∈ *range f*)

A function f is bijective on R^{++} iff it is injective and surjective on R^{++}.

abbreviation *bij-on-trancl R f* ≡ *inj f* ∧ *surj-on-trancl R f*

1.2.2 Helper Lemmas

lemma *rtranclp-eq-rtranclp* [*iff*]:
(λ*x y*. *P x y* ∨ *x* = *y*)** = P^{**}
proof (*intro ext iffI*)
fix *x y*
have (λ*x y*. *P x y* ∨ *x* = *y*)** *x y* ⟶ $P^{==**}$ *x y*
by (*rule mono-rtranclp*) *simp*
thus (λ*x y*. *P x y* ∨ *x* = *y*)** *x y* ⟹ P^{**} *x y*
by *simp*
show P^{**} *x y* ⟹ (λ*x y*. *P x y* ∨ *x* = *y*)** *x y*
by (*metis* (*no-types*, *lifting*) *mono-rtranclp*)
qed

lemma *tranclp-eq-rtranclp* [*iff*]:
(λ*x y*. *P x y* ∨ *x* = *y*)$^{++}$ = P^{**}
proof (*intro ext iffI*)
fix *x y*
have (λ*x y*. *P x y* ∨ *x* = *y*)** *x y* ⟶ $P^{==**}$ *x y*
by (*rule mono-rtranclp*) *simp*
thus (λ*x y*. *P x y* ∨ *x* = *y*)$^{++}$ *x y* ⟹ P^{**} *x y*
using *tranclp-into-rtranclp* **by** *force*
show P^{**} *x y* ⟹ (λ*x y*. *P x y* ∨ *x* = *y*)$^{++}$ *x y*
by (*metis* (*mono-tags*, *lifting*) *mono-rtranclp rtranclpD tranclp.r-into-trancl*)
qed

lemma *rtranclp-eq-rtranclp′* [*iff*]:
(λ*x y*. *P x y* ∧ *x* ≠ *y*)** = P^{**}
proof (*intro ext iffI*)
fix *x y*
show (λ*x y*. *P x y* ∧ *x* ≠ *y*)** *x y* ⟹ P^{**} *x y*
by (*metis* (*no-types*, *lifting*) *mono-rtranclp*)

assume $P^{**}\ x\ y$
hence $(inf\ P\ (\neq))^{**}\ x\ y$
by (*simp add*: *rtranclp-r-diff-Id*)
also have $(inf\ P\ (\neq))^{**}\ x\ y \longrightarrow (\lambda x\ y.\ P\ x\ y \wedge x \neq y)^{**}\ x\ y$
by (*rule mono-rtranclp*) *simp*
finally show $P^{**}\ x\ y \Longrightarrow (\lambda x\ y.\ P\ x\ y \wedge x \neq y)^{**}\ x\ y$ **by** *simp*
qed

lemma *tranclp-tranclp-to-tranclp-r*:
assumes $(\bigwedge x\ y\ z.\ R^{++}\ x\ y \Longrightarrow R\ y\ z \Longrightarrow P\ x \Longrightarrow P\ z \Longrightarrow P\ y)$
assumes $R^{++}\ x\ y$ **and** $R^{++}\ y\ z$
assumes $P\ x$ **and** $P\ z$
shows $P\ y$
proof −
have $(\bigwedge x\ y\ z.\ R^{++}\ x\ y \Longrightarrow R\ y\ z \Longrightarrow P\ x \Longrightarrow P\ z \Longrightarrow P\ y) \Longrightarrow$
$R^{++}\ y\ z \Longrightarrow R^{++}\ x\ y \Longrightarrow P\ x \longrightarrow P\ z \longrightarrow P\ y$
by (*erule tranclp-induct*, *auto*) (*meson tranclp-trans*)
thus *?thesis* **using** *assms* **by** *auto*
qed

1.2.3 Transitive Closure Preservation

A function f preserves R^{++} if it preserves R.

The proof was derived from the accepted answer on the website Stack Overflow that is available at https://stackoverflow.com/a/52573551/632199 and provided with the permission of the author of the answer.

lemma *preserve-tranclp*:
assumes $\bigwedge x\ y.\ R\ x\ y \Longrightarrow S\ (f\ x)\ (f\ y)$
assumes $R^{++}\ x\ y$
shows $S^{++}\ (f\ x)\ (f\ y)$
proof −
define P **where** P: $P = (\lambda x\ y.\ S^{++}\ (f\ x)\ (f\ y))$
define r **where** r: $r = (\lambda x\ y.\ S\ (f\ x)\ (f\ y))$
have $r^{++}\ x\ y$ **by** (*insert assms r*; *erule tranclp-trans-induct*; *auto*)
moreover have $\bigwedge x\ y.\ r\ x\ y \Longrightarrow P\ x\ y$ **unfolding** $P\ r$ **by** *simp*
moreover have $\bigwedge x\ y\ z.\ r^{++}\ x\ y \Longrightarrow P\ x\ y \Longrightarrow r^{++}\ y\ z \Longrightarrow P\ y\ z \Longrightarrow P\ x\ z$
unfolding P **by** *auto*
ultimately have $P\ x\ y$ **by** (*rule tranclp-trans-induct*)
with P **show** *?thesis* **by** *simp*
qed

A function f preserves R^{**} if it preserves R.

lemma *preserve-rtranclp*:
$(\bigwedge x\ y.\ R\ x\ y \Longrightarrow S\ (f\ x)\ (f\ y)) \Longrightarrow$
$R^{**}\ x\ y \Longrightarrow S^{**}\ (f\ x)\ (f\ y)$
unfolding *Nitpick.rtranclp-unfold*
by (*metis preserve-tranclp*)

If one needs to prove that $(f\ x)$ and $(g\ y)$ are related through S^{**} then one can use the previous lemma and add a one more step from $(f\ y)$ to $(g\ y)$.

lemma *preserve-rtranclp'*:
$(\bigwedge x\ y.\ R\ x\ y \Longrightarrow S\ (f\ x)\ (f\ y)) \Longrightarrow$
$(\bigwedge y.\ S\ (f\ y)\ (g\ y)) \Longrightarrow$
$R^{**}\ x\ y \Longrightarrow S^{**}\ (f\ x)\ (g\ y)$
by (*metis preserve-rtranclp rtranclp.rtrancl-into-rtrancl*)

lemma *preserve-rtranclp''*:
$(\bigwedge x\ y.\ R\ x\ y \Longrightarrow S\ (f\ x)\ (f\ y)) \Longrightarrow$
$(\bigwedge y.\ S\ (f\ y)\ (g\ y)) \Longrightarrow$
$R^{**}\ x\ y \Longrightarrow S^{++}\ (f\ x)\ (g\ y)$
apply (*rule-tac ?b= f y* **in** *rtranclp-into-tranclp1*, *auto*)
by (*rule preserve-rtranclp*, *auto*)

1.2.4 Transitive Closure Reflection

A function f reflects S^{++} if it reflects S and is bijective on S^{++}.

The proof was derived from the accepted answer on the website Stack Overflow that is available at https://stackoverflow.com/a/52573551/632199 and provided with the permission of the author of the answer.

lemma *reflect-tranclp*:
assumes *refl-f*: $\bigwedge x\ y.\ S\ (f\ x)\ (f\ y) \Longrightarrow R\ x\ y$
assumes *bij-f*: *bij-on-trancl S f*
assumes *prem*: $S^{++}\ (f\ x)\ (f\ y)$
shows $R^{++}\ x\ y$
proof −
define B **where** B: $B = range\ f$
define g **where** g: $g = the\text{-}inv\text{-}into\ UNIV\ f$
define gr **where** gr: $gr = restrict\ g\ B$
define P **where** P: $P = (\lambda x\ y.\ x \in B \longrightarrow y \in B \longrightarrow R^{++}\ (gr\ x)\ (gr\ y))$
from *prem* **have** *major*: $S^{++}\ (f\ x)\ (f\ y)$ **by** *blast*
from *refl-f bij-f* **have** *cases-1*: $\bigwedge x\ y.\ S\ x\ y \Longrightarrow P\ x\ y$
unfolding *B P g gr*
by (*simp add*: *f-the-inv-into-f tranclp.r-into-trancl*)
from *refl-f bij-f*
have $(\bigwedge x\ y\ z.\ S^{++}\ x\ y \Longrightarrow S^{++}\ y\ z \Longrightarrow x \in B \Longrightarrow z \in B \Longrightarrow y \in B)$
unfolding *B*
by (*rule-tac ?z= z* **in** *tranclp-tranclp-to-tranclp-r*, *auto*, *blast*)
with *P* **have** *cases-2*:
$\bigwedge x\ y\ z.\ S^{++}\ x\ y \Longrightarrow P\ x\ y \Longrightarrow S^{++}\ y\ z \Longrightarrow P\ y\ z \Longrightarrow P\ x\ z$
unfolding *B*
by *auto*
from *major cases-1 cases-2* **have** $P\ (f\ x)\ (f\ y)$
by (*rule tranclp-trans-induct*)
with *bij-f* **show** *?thesis* **unfolding** *P B g gr* **by** (*simp add*: *the-inv-f-f*)

qed

A function f reflects S^{**} if it reflects S and is bijective on S^{++}.

lemma *reflect-rtranclp*:
$(\bigwedge x\ y.\ S\ (f\ x)\ (f\ y) \Longrightarrow R\ x\ y) \Longrightarrow$
bij-on-trancl $S\ f \Longrightarrow$
$S^{**}\ (f\ x)\ (f\ y) \Longrightarrow R^{**}\ x\ y$
unfolding *Nitpick.rtranclp-unfold*
by (*metis* (*full-types*) *injD reflect-tranclp*)

end

1.3 Finite Maps

theory *Finite-Map-Ext*
imports *HOL−Library.Finite-Map*
begin

type-notation *fmap* (‹(- $\rightharpoonup_f$ /-)› [*22*, *21*] *21*)

nonterminal *fmaplets* **and** *fmaplet*

syntax
-fmaplet :: ['*a*, '*a*] ⇒ *fmaplet* (‹- /↦$_f$/ -›)
-fmaplets :: ['*a*, '*a*] ⇒ *fmaplet* (‹- /[↦$_f$]/ -›)
:: *fmaplet* ⇒ *fmaplets* (‹-›)
-FMaplets :: [*fmaplet*, *fmaplets*] ⇒ *fmaplets* (‹-,/ -›)
-FMapUpd :: ['*a* ⇀ '*b*, *fmaplets*] ⇒ '*a* ⇀ '*b* (‹-/'(-')› [*900*, *0*] *900*)
-FMap :: *fmaplets* ⇒ '*a* ⇀ '*b* (‹(*1*[-])›)

syntax (*ASCII*)
-fmaplet :: ['*a*, '*a*] ⇒ *fmaplet* (‹- /|->f/ -›)
-fmaplets :: ['*a*, '*a*] ⇒ *fmaplet* (‹- /[|->f]/ -›)

syntax-consts
-fmaplet *-fmaplets* *-FMaplets* *-FMapUpd* *-FMap* ⇌ *fmupd*

translations
-FMapUpd m (*-FMaplets xy ms*) ⇌ *-FMapUpd* (*-FMapUpd m xy*) *ms*
-FMapUpd m (*-fmaplet x y*) ⇌ *CONST fmupd x y m*
-FMap ms ⇌ *-FMapUpd* (*CONST fmempty*) *ms*
-FMap (*-FMaplets ms1 ms2*) ↽ *-FMapUpd* (*-FMap ms1*) *ms2*
-FMaplets ms1 (*-FMaplets ms2 ms3*) ↽ *-FMaplets* (*-FMaplets ms1 ms2*) *ms3*

1.3.1 Helper Lemmas

lemma *fmrel-on-fset-fmdom*:
fmrel-on-fset (*fmdom ym*) *f xm ym* ⟹
k |∈| *fmdom ym* ⟹

k |∈| *fmdom xm*
by (*metis fmdom-notD fmdom-notI fmrel-on-fsetD option.rel-sel*)

1.3.2 Merge Operation

definition *fmmerge f xm ym* ≡
fmap-of-list (*map*
(λ*k*. (*k*, *f* (*the* (*fmlookup xm k*)) (*the* (*fmlookup ym k*))))
(*sorted-list-of-fset* (*fmdom xm* |∩| *fmdom ym*)))

lemma *fmdom-fmmerge* [*simp*]:
fmdom (*fmmerge g xm ym*) = *fmdom xm* |∩| *fmdom ym*
by (*auto simp add*: *fmmerge-def fmdom-of-list*)

lemma *fmmerge-commut*:
assumes ⋀*x y*. *x* ∈ *fmran′ xm* ⟹ *f x y* = *f y x*
shows *fmmerge f xm ym* = *fmmerge f ym xm*
proof −
obtain *zm* **where** *zm*: *zm* = *sorted-list-of-fset* (*fmdom xm* |∩| *fmdom ym*)
by *auto*
with *assms* **have**
map (λ*k*. (*k*, *f* (*the* (*fmlookup xm k*)) (*the* (*fmlookup ym k*)))) *zm* =
map (λ*k*. (*k*, *f* (*the* (*fmlookup ym k*)) (*the* (*fmlookup xm k*)))) *zm*
by (*auto*) (*metis fmdom-notI fmran′I option.collapse*)
thus *?thesis*
unfolding *fmmerge-def zm*
by (*metis* (*no-types*, *lifting*) *inf-commute*)
qed

lemma *fmrel-on-fset-fmmerge1* [*intro*]:
assumes ⋀*x y z*. *z* ∈ *fmran′ zm* ⟹ *f x z* ⟹ *f y z* ⟹ *f* (*g x y*) *z*
assumes *fmrel-on-fset* (*fmdom zm*) *f xm zm*
assumes *fmrel-on-fset* (*fmdom zm*) *f ym zm*
shows *fmrel-on-fset* (*fmdom zm*) *f* (*fmmerge g xm ym*) *zm*
proof −
{
fix *x a b c*
assume *x* |∈| *fmdom zm*
moreover hence *x* |∈| *fmdom xm* |∩| *fmdom ym*
by (*meson assms*(*2*) *assms*(*3*) *finterI fmrel-on-fset-fmdom*)
moreover assume *fmlookup xm x* = *Some a*
and *fmlookup ym x* = *Some b*
and *fmlookup zm x* = *Some c*
moreover from *assms calculation* **have** *f* (*g a b*) *c*
by (*metis fmran′I fmrel-on-fsetD option.rel-inject*(*2*))
ultimately have
rel-option f (*fmlookup* (*fmmerge g xm ym*) *x*) (*fmlookup zm x*)
unfolding *fmmerge-def fmlookup-of-list* **apply** *auto*
unfolding *option-rel-Some2* **apply** (*rule-tac ?x*= *g a b* **in** *exI*)

unfolding *map-of-map-restrict restrict-map-def*
by *auto*
}
with *assms(2) assms(3)* **show** *?thesis*
by (*meson fmdomE fmrel-on-fsetI fmrel-on-fset-fmdom*)
qed

lemma *fmrel-on-fset-fmmerge2* [*intro*]:
assumes $\bigwedge x\ y.\ x \in fmran'\ xm \Longrightarrow f\ x\ (g\ x\ y)$
shows *fmrel-on-fset* (*fmdom ym*) *f xm* (*fmmerge g xm ym*)
proof −
{
fix *x a b*
assume $x \mid\in\mid fmdom\ xm \mid\cap\mid fmdom\ ym$
and *fmlookup xm x* = *Some a*
and *fmlookup ym x* = *Some b*
hence *rel-option f* (*fmlookup xm x*) (*fmlookup* (*fmmerge g xm ym*) *x*)
unfolding *fmmerge-def fmlookup-of-list* **apply** *auto*
unfolding *option-rel-Some1* **apply** (*rule-tac ?x*= *g a b* **in** *exI*)
by (*auto simp add*: *map-of-map-restrict assms fmran'I*)
}
with *assms* **show** *?thesis*
apply *auto*
apply (*rule fmrel-on-fsetI*)
by (*metis* (*full-types*) *finterD1 fmdomE fmdom-fmmerge fmdom-notD rel-option-None2*)
qed

1.3.3 Acyclicity

abbreviation *acyclic-on xs r* ≡ $(\forall x.\ x \in xs \longrightarrow (x,\ x) \notin r^+)$

abbreviation *acyclicP-on xs r* ≡ *acyclic-on xs* $\{(x,\ y).\ r\ x\ y\}$

lemma *fmrel-acyclic*:
acyclicP-on (*fmran' xm*) *R* ⟹
fmrel R^{++} *xm ym* ⟹
fmrel R ym xm ⟹
xm = *ym*
by (*metis* (*full-types*) *fmap-ext fmran'I fmrel-cases option.sel*
tranclp.trancl-into-trancl tranclp-unfold)

lemma *fmrel-acyclic'*:
assumes *acyclicP-on* (*fmran' ym*) *R*
assumes *fmrel* R^{++} *xm ym*
assumes *fmrel R ym xm*
shows *xm* = *ym*
proof −
{
fix *x*

```
    from assms(1) have
      rel-option R^++ (fmlookup xm x) (fmlookup ym x) ⟹
      rel-option R (fmlookup ym x) (fmlookup xm x) ⟹
      rel-option R (fmlookup xm x) (fmlookup ym x)
     by (metis (full-types) fmdom'-notD fmlookup-dom'-iff
          fmran'I option.rel-sel option.sel
          tranclp-into-tranclp2 tranclp-unfold)
  }
  with assms show ?thesis
   unfolding fmrel-iff
   by (metis fmap.rel-mono-strong fmrelI fmrel-acyclic tranclp.simps)
qed

lemma fmrel-on-fset-acyclic:
  acyclicP-on (fmran' xm) R ⟹
  fmrel-on-fset (fmdom ym) R^++ xm ym ⟹
  fmrel-on-fset (fmdom xm) R ym xm ⟹
  xm = ym
 unfolding fmrel-on-fset-fmrel-restrict
 by (metis (no-types, lifting) fmdom-filter fmfilter-alt-defs(5)
      fmfilter-cong fmlookup-filter fmrel-acyclic fmrel-fmdom-eq
      fmrestrict-fset-dom option.simps(3))

lemma fmrel-on-fset-acyclic':
  acyclicP-on (fmran' ym) R ⟹
  fmrel-on-fset (fmdom ym) R^++ xm ym ⟹
  fmrel-on-fset (fmdom xm) R ym xm ⟹
  xm = ym
 unfolding fmrel-on-fset-fmrel-restrict
 by (metis (no-types, lifting) ffmember-filter fmdom-filter
      fmfilter-alt-defs(5) fmfilter-cong fmrel-acyclic'
      fmrel-fmdom-eq fmrestrict-fset-dom)
```

1.3.4 Transitive Closures

```
lemma fmrel-trans:
  (⋀x y z. x ∈ fmran' xm ⟹ P x y ⟹ Q y z ⟹ R x z) ⟹
  fmrel P xm ym ⟹ fmrel Q ym zm ⟹ fmrel R xm zm
 unfolding fmrel-iff
 by (metis fmdomE fmdom-notD fmran'I option.rel-inject(2) option.rel-sel)

lemma fmrel-on-fset-trans:
  (⋀x y z. x ∈ fmran' xm ⟹ P x y ⟹ Q y z ⟹ R x z) ⟹
  fmrel-on-fset (fmdom ym) P xm ym ⟹
  fmrel-on-fset (fmdom zm) Q ym zm ⟹
  fmrel-on-fset (fmdom zm) R xm zm
 apply (rule fmrel-on-fsetI)
 unfolding option.rel-sel apply auto
 apply (meson fmdom-notI fmrel-on-fset-fmdom)
```

by (*metis fmdom-notI fmran'I fmrel-on-fsetD fmrel-on-fset-fmdom option.rel-sel option.sel*)

lemma *trancl-to-fmrel*:
(*fmrel f*)$^{++}$ *xm ym* $\Longrightarrow$ *fmrel f*$^{++}$ *xm ym*
apply (*induct rule*: *tranclp-induct*)
apply (*simp add*: *fmap.rel-mono-strong*)
by (*rule fmrel-trans*; *auto*)

lemma *fmrel-trancl-fmdom-eq*:
(*fmrel f*)$^{++}$ *xm ym* $\Longrightarrow$ *fmdom xm* = *fmdom ym*
by (*induct rule*: *tranclp-induct*; *simp add*: *fmrel-fmdom-eq*)

The proof was derived from the accepted answer on the website Stack Overflow that is available at https://stackoverflow.com/a/53585232/632199 and provided with the permission of the author of the answer.

lemma *fmupd-fmdrop*:
fmlookup xm k = *Some x* $\Longrightarrow$
xm = *fmupd k x* (*fmdrop k xm*)
apply (*rule fmap-ext*)
unfolding *fmlookup-drop fmupd-lookup*
by *auto*

lemma *fmap-eqdom-Cons1*:
assumes *fmlookup xm i* = *None*
and *fmdom* (*fmupd i x xm*) = *fmdom ym*
and *fmrel R* (*fmupd i x xm*) *ym*
shows ($\exists$ *z zm*. *fmlookup zm i* = *None* $\wedge$ *ym* = (*fmupd i z zm*) $\wedge$
R x z $\wedge$ *fmrel R xm zm*)
proof −
from *assms*(*2*) **obtain** *y* **where** *fmlookup ym i* = *Some y* **by** *force*
then obtain *z zm* **where** *z-zm*: *ym* = *fmupd i z zm* $\wedge$ *fmlookup zm i* = *None*
using *fmupd-fmdrop* **by** *force*
{
assume ¬ *R x z*
with *z-zm* **have** ¬ *fmrel R* (*fmupd i x xm*) *ym*
by (*metis fmrel-iff fmupd-lookup option.simps*(*11*))
}
with *assms*(*3*) **moreover have** *R x z* **by** *auto*
{
assume ¬ *fmrel R xm zm*
with *assms*(*1*) **have** ¬ *fmrel R* (*fmupd i x xm*) *ym*
by (*metis fmrel-iff fmupd-lookup option.rel-sel z-zm*)
}
with *assms*(*3*) **moreover have** *fmrel R xm zm* **by** *auto*
ultimately show *?thesis* **using** *z-zm* **by** *blast*
qed

The proof was derived from the accepted answer on the website Stack Overflow that is available at https://stackoverflow.com/a/53585232/632199

and provided with the permission of the author of the answer.

```
lemma fmap-eqdom-induct [consumes 2, case-names nil step]:
  assumes R: fmrel R xm ym
    and dom-eq: fmdom xm = fmdom ym
    and nil: P (fmap-of-list []) (fmap-of-list [])
    and step:
     ⋀x xm y ym i.
    ⟦R x y; fmrel R xm ym; fmdom xm = fmdom ym; P xm ym⟧ ⟹
    P (fmupd i x xm) (fmupd i y ym)
  shows P xm ym
  using R dom-eq
proof (induct xm arbitrary: ym)
  case fmempty thus ?case
    by (metis fempty-iff fmdom-empty fmempty-of-list fmfilter-alt-defs(5)
          fmfilter-false fmrestrict-fset-dom local.nil)
next
  case (fmupd i x xm) show ?case
  proof −
    obtain y where fmlookup ym i = Some y
      by (metis fmupd.prems(1) fmrel-cases fmupd-lookup option.discI)
    from fmupd.hyps(2) fmupd.prems(1) fmupd.prems(2) obtain z zm where
      fmlookup zm i = None and
      ym-eq-z-zm: ym = (fmupd i z zm) and
      R-x-z: R x z and
      R-xm-zm: fmrel R xm zm
      using fmap-eqdom-Cons1 by metis
    hence dom-xm-eq-dom-zm: fmdom xm = fmdom zm
      using fmrel-fmdom-eq by blast
    with R-xm-zm fmupd.hyps(1) have P xm zm by blast
    with R-x-z R-xm-zm dom-xm-eq-dom-zm have
      P (fmupd i x xm) (fmupd i z zm)
      by (rule step)
    thus ?thesis by (simp add: ym-eq-z-zm)
  qed
qed
```

The proof was derived from the accepted answer on the website Stack Overflow that is available at https://stackoverflow.com/a/53585232/632199 and provided with the permission of the author of the answer.

```
lemma fmrel-to-rtrancl:
  assumes as-r: reflp r
     and rel-rpp-xm-ym: fmrel r** xm ym
   shows (fmrel r)** xm ym
proof −
  from rel-rpp-xm-ym have fmdom xm = fmdom ym
    using fmrel-fmdom-eq by blast
  with rel-rpp-xm-ym show (fmrel r)** xm ym
  proof (induct rule: fmap-eqdom-induct)
    case nil show ?case by auto
```

```
  next
    case (step x xm y ym i) show ?case
    proof −
      from step.hyps(1) have (fmrel r)** (fmupd i x xm) (fmupd i y xm)
      proof (induct rule: rtranclp-induct)
        case base show ?case by simp
      next
        case (step y z) show ?case
        proof −
          from as-r have fmrel r xm xm
            by (simp add: fmap.rel-reflp reflpD)
          with step.hyps(2) have (fmrel r)** (fmupd i y xm) (fmupd i z xm)
            by (simp add: fmrel-upd r-into-rtranclp)
          with step.hyps(3) show ?thesis by simp
        qed
      qed
      also from step.hyps(4) have (fmrel r)** (fmupd i y xm) (fmupd i y ym)
      proof (induct rule: rtranclp-induct)
        case base show ?case by simp
      next
        case (step ya za) show ?case
        proof −
          from step.hyps(2) as-r have (fmrel r)** (fmupd i y ya) (fmupd i y za)
            by (simp add: fmrel-upd r-into-rtranclp reflp-def)
          with step.hyps(3) show ?thesis by simp
        qed
      qed
      finally show ?thesis by simp
    qed
  qed
qed
```

The proof was derived from the accepted answer on the website Stack Overflow that is available at https://stackoverflow.com/a/53585232/632199 and provided with the permission of the author of the answer.

```
lemma fmrel-to-trancl:
  assumes reflp r
      and fmrel r++ xm ym
    shows (fmrel r)++ xm ym
proof −
  from assms(2) have fmrel r** xm ym
    by (drule-tac ?Ra= r** in fmap.rel-mono-strong; auto)
  with assms(1) have (fmrel r)** xm ym
    by (simp add: fmrel-to-rtrancl)
  with assms(1) show ?thesis
    by (metis fmap.rel-reflp reflpD rtranclpD tranclp.r-into-trancl)
qed

lemma fmrel-tranclp-induct:
```

fmrel r^{++} *a b* $\Longrightarrow$
reflp r $\Longrightarrow$
$(\bigwedge y.$ *fmrel r a y* $\Longrightarrow$ *P y*$) \Longrightarrow$
$(\bigwedge y\ z.$ $(fmrel\ r)^{++}$ *a y* $\Longrightarrow$ *fmrel r y z* $\Longrightarrow$ *P y* $\Longrightarrow$ *P z*$) \Longrightarrow$ *P b*
apply (*drule fmrel-to-trancl*, *simp*)
by (*erule tranclp-induct*; *simp*)

lemma *fmrel-converse-tranclp-induct*:
fmrel r^{++} *a b* $\Longrightarrow$
reflp r $\Longrightarrow$
$(\bigwedge y.$ *fmrel r y b* $\Longrightarrow$ *P y*$) \Longrightarrow$
$(\bigwedge y\ z.$ *fmrel r y z* $\Longrightarrow$ *fmrel* r^{++} *z b* $\Longrightarrow$ *P z* $\Longrightarrow$ *P y*$) \Longrightarrow$ *P a*
apply (*drule fmrel-to-trancl*, *simp*)
by (*erule converse-tranclp-induct*; *simp add*: *trancl-to-fmrel*)

lemma *fmrel-tranclp-trans-induct*:
fmrel r^{++} *a b* $\Longrightarrow$
reflp r $\Longrightarrow$
$(\bigwedge x\ y.$ *fmrel r x y* $\Longrightarrow$ *P x y*$) \Longrightarrow$
$(\bigwedge x\ y\ z.$ *fmrel* r^{++} *x y* $\Longrightarrow$ *P x y* $\Longrightarrow$ *fmrel* r^{++} *y z* $\Longrightarrow$ *P y z* $\Longrightarrow$ *P x z*$) \Longrightarrow$
P a b
apply (*drule fmrel-to-trancl*, *simp*)
apply (*erule tranclp-trans-induct*, *simp*)
using *trancl-to-fmrel* **by** *blast*

1.3.5 Size Calculation

The contents of the subsection was derived from the accepted answer on the website Stack Overflow that is available at https://stackoverflow.com/a/53244203/632199 and provided with the permission of the author of the answer.

abbreviation *tcf* $\equiv (\lambda\ v{::}('a \times nat).\ (\lambda\ r{::}nat.\ snd\ v + r))$

interpretation *tcf*: *comp-fun-commute tcf*
proof
fix *x y* :: $'a \times nat$
show *tcf y* $\circ$ *tcf x* $=$ *tcf x* $\circ$ *tcf y*
proof −
fix *z*
have (*tcf y* $\circ$ *tcf x*) *z* $=$ *snd y* $+$ *snd x* $+$ *z* **by** *auto*
also have (*tcf x* $\circ$ *tcf y*) *z* $=$ *snd y* $+$ *snd x* $+$ *z* **by** *auto*
finally have (*tcf y* $\circ$ *tcf x*) *z* $=$ (*tcf x* $\circ$ *tcf y*) *z* **by** *auto*
thus (*tcf y* $\circ$ *tcf x*) $=$ (*tcf x* $\circ$ *tcf y*) **by** *auto*
qed
qed

lemma *ffold-rec-exp*:
assumes *k* $|\in|$ *fmdom x*
and *ky* $=$ (*k*, *the* (*fmlookup* (*fmmap f x*) *k*))

shows *ffold tcf 0 (fset-of-fmap (fmmap f x)) =*
tcf ky (ffold tcf 0 ((fset-of-fmap (fmmap f x)) |−| {|ky|}))
proof −
have *ky |∈| (fset-of-fmap (fmmap f x))*
using *assms* **by** *auto*
thus *?thesis*
by (*simp add*: *tcf.ffold-rec*)
qed

lemma *elem-le-ffold* [*intro*]:
k |∈| fmdom x ⟹
f (the (fmlookup x k)) < Suc (ffold tcf 0 (fset-of-fmap (fmmap f x)))
by (*subst ffold-rec-exp*, *auto*)

lemma *elem-le-ffold′* [*intro*]:
z ∈ fmran′ x ⟹
f z < Suc (ffold tcf 0 (fset-of-fmap (fmmap f x)))
apply (*erule fmran′E*)
apply (*frule fmdomI*)
by (*subst ffold-rec-exp*, *auto*)

1.3.6 Code Setup

abbreviation *fmmerge-fun f xm ym ≡*
fmap-of-list (map
(λk. if k |∈| fmdom xm ∧ k |∈| fmdom ym
then (k, f (the (fmlookup xm k)) (the (fmlookup ym k)))
else (k, undefined))
(sorted-list-of-fset (fmdom xm |∩| fmdom ym)))

lemma *fmmerge-fun-simp* [*code-abbrev*, *simp*]:
fmmerge-fun f xm ym = fmmerge f xm ym
unfolding *fmmerge-def*
apply (*rule-tac ?f= fmap-of-list* **in** *HOL.arg-cong*)
by *simp*

end

1.4 Tuples

theory *Tuple*
imports *Finite-Map-Ext Transitive-Closure-Ext*
begin

1.4.1 Definitions

abbreviation *subtuple f xm ym ≡ fmrel-on-fset (fmdom ym) f xm ym*

abbreviation *strict-subtuple f xm ym ≡ subtuple f xm ym ∧ xm ≠ ym*

1.4.2 Helper Lemmas

lemma *fmrel-to-subtuple*:
fmrel r xm ym $\Longrightarrow$ *subtuple r xm ym*
unfolding *fmrel-on-fset-fmrel-restrict* **by** *blast*

lemma *subtuple-eq-fmrel-fmrestrict-fset*:
subtuple r xm ym = *fmrel r* (*fmrestrict-fset* (*fmdom ym*) *xm*) *ym*
by (*simp add*: *fmrel-on-fset-fmrel-restrict*)

lemma *subtuple-fmdom*:
subtuple f xm ym $\Longrightarrow$
subtuple g ym xm $\Longrightarrow$
fmdom xm = *fmdom ym*
by (*meson fmrel-on-fset-fmdom fset-eqI*)

1.4.3 Basic Properties

lemma *subtuple-refl*:
reflp R $\Longrightarrow$ *subtuple R xm xm*
by (*simp add*: *fmrel-on-fsetI option.rel-reflp reflpD*)

lemma *subtuple-mono* [*mono*]:
($\bigwedge$*x y*. *x* $\in$ *fmran' xm* $\Longrightarrow$ *y* $\in$ *fmran' ym* $\Longrightarrow$ *f x y* $\longrightarrow$ *g x y*) $\Longrightarrow$
subtuple f xm ym $\longrightarrow$ *subtuple g xm ym*
apply (*auto*)
apply (*rule fmrel-on-fsetI*)
apply (*drule-tac ?P*= *f* **and** *?m*= *xm* **and** *?n*= *ym* **in** *fmrel-on-fsetD*, *simp*)
apply (*erule option.rel-cases*, *simp*)
by (*auto simp add*: *option.rel-sel fmran'I*)

lemma *strict-subtuple-mono* [*mono*]:
($\bigwedge$*x y*. *x* $\in$ *fmran' xm* $\Longrightarrow$ *y* $\in$ *fmran' ym* $\Longrightarrow$ *f x y* $\longrightarrow$ *g x y*) $\Longrightarrow$
strict-subtuple f xm ym $\longrightarrow$ *strict-subtuple g xm ym*
using *subtuple-mono* **by** *blast*

lemma *subtuple-antisym*:
assumes *subtuple* (λ*x y*. *f x y* $\vee$ *x* = *y*) *xm ym*
assumes *subtuple* (λ*x y*. *f x y* $\wedge$ $\neg$ *f y x* $\vee$ *x* = *y*) *ym xm*
shows *xm* = *ym*
proof (*rule fmap-ext*)
fix *x*
from *assms* **have** *fmdom xm* = *fmdom ym*
using *subtuple-fmdom* **by** *blast*
with *assms* **have** *fmrel* (λ*x y*. *f x y* $\vee$ *x* = *y*) *xm ym*
and *fmrel* (λ*x y*. *f x y* $\wedge$ $\neg$ *f y x* $\vee$ *x* = *y*) *ym xm*
by (*metis* (*mono-tags*, *lifting*) *fmrel-code fmrel-on-fset-alt-def*)+
thus *fmlookup xm x* = *fmlookup ym x*
apply (*erule-tac ?x*= *x* **in** *fmrel-cases*)
by (*erule-tac ?x*= *x* **in** *fmrel-cases*, *auto*)+

qed

lemma *strict-subtuple-antisym*:
  strict-subtuple $(\lambda x\ y.\ f\ x\ y \vee x = y)$ *xm ym* $\Longrightarrow$
  strict-subtuple $(\lambda x\ y.\ f\ x\ y \wedge \neg\ f\ y\ x \vee x = y)$ *ym xm* $\Longrightarrow$ *False*
 by (*auto simp add*: *subtuple-antisym*)

lemma *subtuple-acyclic*:
 assumes *acyclicP-on* (*fmran'* *xm*) *P*
 assumes *subtuple* $(\lambda x\ y.\ P\ x\ y \vee x = y)^{++}$ *xm ym*
 assumes *subtuple* $(\lambda x\ y.\ P\ x\ y \vee x = y)$ *ym xm*
 shows *xm* = *ym*
proof (*rule fmap-ext*)
 fix *x*
 from *assms* **have** *fmdom-eq*: *fmdom xm* = *fmdom ym*
  using *subtuple-fmdom* **by** *blast*
 have $\bigwedge x\ a\ b.$ *acyclicP-on* (*fmran'* *xm*) *P* $\Longrightarrow$
   fmlookup xm x = *Some a* $\Longrightarrow$
   fmlookup ym x = *Some b* $\Longrightarrow$
   $P^{**}\ a\ b \Longrightarrow P\ b\ a \vee a = b \Longrightarrow a = b$
  by (*meson Nitpick.tranclp-unfold fmran'I rtranclp-into-tranclp1*)
 moreover from *fmdom-eq assms*(*2*) **have** *fmrel* P^{**} *xm ym*
  unfolding *fmrel-on-fset-fmrel-restrict* **apply** *auto*
  by (*metis fmrestrict-fset-dom*)
 moreover from *fmdom-eq assms*(*3*) **have** *fmrel* $(\lambda x\ y.\ P\ x\ y \vee x = y)$ *ym xm*
  unfolding *fmrel-on-fset-fmrel-restrict* **apply** *auto*
  by (*metis fmrestrict-fset-dom*)
 ultimately show *fmlookup xm x* = *fmlookup ym x*
  apply (*erule-tac ?x*= *x* **in** *fmrel-cases*)
  apply (*erule-tac ?x*= *x* **in** *fmrel-cases*, *simp-all*)+
  using *assms*(*1*) **by** *blast*
qed

lemma *subtuple-acyclic'*:
 assumes *acyclicP-on* (*fmran'* *ym*) *P*
 assumes *subtuple* $(\lambda x\ y.\ P\ x\ y \vee x = y)^{++}$ *xm ym*
 assumes *subtuple* $(\lambda x\ y.\ P\ x\ y \vee x = y)$ *ym xm*
 shows *xm* = *ym*
proof (*rule fmap-ext*)
 fix *x*
 from *assms* **have** *fmdom-eq*: *fmdom xm* = *fmdom ym*
  using *subtuple-fmdom* **by** *blast*
 have $\bigwedge x\ a\ b.$ *acyclicP-on* (*fmran'* *ym*) *P* $\Longrightarrow$
   fmlookup xm x = *Some a* $\Longrightarrow$
   fmlookup ym x = *Some b* $\Longrightarrow$
   $P^{**}\ a\ b \Longrightarrow P\ b\ a \vee a = b \Longrightarrow a = b$
  by (*meson Nitpick.tranclp-unfold fmran'I rtranclp-into-tranclp2*)
 moreover from *fmdom-eq assms*(*2*) **have** *fmrel* P^{**} *xm ym*
  unfolding *fmrel-on-fset-fmrel-restrict* **apply** *auto*

by (*metis fmrestrict-fset-dom*)
moreover from *fmdom-eq assms*(*3*) **have** *fmrel* (λ*x y*. *P x y* ∨ *x* = *y*) *ym xm*
unfolding *fmrel-on-fset-fmrel-restrict* **apply** *auto*
by (*metis fmrestrict-fset-dom*)
ultimately show *fmlookup xm x* = *fmlookup ym x*
apply (*erule-tac ?x*= *x* **in** *fmrel-cases*)
apply (*erule-tac ?x*= *x* **in** *fmrel-cases*, *simp-all*)+
using *assms*(*1*) **by** *blast*
qed

lemma *subtuple-acyclic″*:
acyclicP-on (*fmran′ ym*) *R* ⟹
*subtuple R*** *xm ym* ⟹
subtuple R ym xm ⟹
xm = *ym*
by (*metis* (*no-types*, *lifting*) *subtuple-acyclic′ subtuple-mono tranclp-eq-rtranclp*)

lemma *strict-subtuple-trans*:
acyclicP-on (*fmran′ xm*) *P* ⟹
strict-subtuple (λ*x y*. *P x y* ∨ *x* = *y*)^++ *xm ym* ⟹
strict-subtuple (λ*x y*. *P x y* ∨ *x* = *y*) *ym zm* ⟹
strict-subtuple (λ*x y*. *P x y* ∨ *x* = *y*)^++ *xm zm*
apply *auto*
apply (*rule fmrel-on-fset-trans*, *auto*)
by (*drule-tac ?ym*= *ym* **in** *subtuple-acyclic*; *auto*)

lemma *strict-subtuple-trans′*:
acyclicP-on (*fmran′ zm*) *P* ⟹
strict-subtuple (λ*x y*. *P x y* ∨ *x* = *y*) *xm ym* ⟹
strict-subtuple (λ*x y*. *P x y* ∨ *x* = *y*)^++ *ym zm* ⟹
strict-subtuple (λ*x y*. *P x y* ∨ *x* = *y*)^++ *xm zm*
apply *auto*
apply (*rule fmrel-on-fset-trans*, *auto*)
by (*drule-tac ?xm*= *ym* **in** *subtuple-acyclic′*; *auto*)

lemma *strict-subtuple-trans″*:
acyclicP-on (*fmran′ zm*) *R* ⟹
strict-subtuple R xm ym ⟹
*strict-subtuple R*** *ym zm* ⟹
*strict-subtuple R*** *xm zm*
apply *auto*
apply (*rule fmrel-on-fset-trans*, *auto*)
by (*drule-tac ?xm*= *ym* **in** *subtuple-acyclic″*; *auto*)

lemma *strict-subtuple-trans‴*:
acyclicP-on (*fmran′ zm*) *P* ⟹
strict-subtuple (λ*x y*. *P x y* ∨ *x* = *y*) *xm ym* ⟹
strict-subtuple (λ*x y*. *P x y* ∨ *x* = *y*)** *ym zm* ⟹
strict-subtuple (λ*x y*. *P x y* ∨ *x* = *y*)** *xm zm*

apply *auto*
apply (*rule fmrel-on-fset-trans*, *auto*)
by (*drule-tac ?xm= ym* **in** *subtuple-acyclic'*; *auto*)

lemma *subtuple-fmmerge2* [*intro*]:
(⋀*x y*. *x* ∈ *fmran' xm* ⟹ *f x* (*g x y*)) ⟹
subtuple f xm (*fmmerge g xm ym*)
by (*rule-tac ?S= fmdom ym* **in** *fmrel-on-fsubset*; *auto*)

1.4.4 Transitive Closures

lemma *trancl-to-subtuple*:
(*subtuple r*)$^{++}$ *xm ym* ⟹
subtuple r^{++} *xm ym*
apply (*induct rule*: *tranclp-induct*)
apply (*metis subtuple-mono tranclp.r-into-trancl*)
by (*rule fmrel-on-fset-trans*; *auto*)

lemma *rtrancl-to-subtuple*:
(*subtuple r*)** *xm ym* ⟹
subtuple r** *xm ym*
apply (*induct rule*: *rtranclp-induct*)
apply (*simp add*: *fmap.rel-refl-strong fmrel-to-subtuple*)
by (*rule fmrel-on-fset-trans*; *auto*)

lemma *fmrel-to-subtuple-trancl*:
reflp r ⟹
(*fmrel r*)$^{++}$ (*fmrestrict-fset* (*fmdom ym*) *xm*) *ym* ⟹
(*subtuple r*)$^{++}$ *xm ym*
apply (*frule trancl-to-fmrel*)
apply (*rule-tac ?r= r* **in** *fmrel-tranclp-induct*, *auto*)
apply (*metis* (*no-types*, *lifting*) *fmrel-fmdom-eq*
subtuple-eq-fmrel-fmrestrict-fset tranclp.r-into-trancl)
by (*meson fmrel-to-subtuple tranclp.simps*)

lemma *subtuple-to-trancl*:
reflp r ⟹
subtuple r^{++} *xm ym* ⟹
(*subtuple r*)$^{++}$ *xm ym*
apply (*rule fmrel-to-subtuple-trancl*)
unfolding *fmrel-on-fset-fmrel-restrict*
by (*simp-all add*: *fmrel-to-trancl*)

lemma *trancl-to-strict-subtuple*:
acyclicP-on (*fmran' ym*) *R* ⟹
(*strict-subtuple R*)$^{++}$ *xm ym* ⟹
strict-subtuple R** *xm ym*
apply (*erule converse-tranclp-induct*)
apply (*metis r-into-rtranclp strict-subtuple-mono*)

using *strict-subtuple-trans''* **by** *blast*

lemma *trancl-to-strict-subtuple'*:
acyclicP-on (*fmran' ym*) *R* ⟹
(*strict-subtuple* (λ*x y*. *R x y* ∨ *x* = *y*))$^{++}$ *xm ym* ⟹
strict-subtuple (λ*x y*. *R x y* ∨ *x* = *y*)** *xm ym*
apply (*erule converse-tranclp-induct*)
apply (*metis* (*no-types*, *lifting*) *r-into-rtranclp strict-subtuple-mono*)
using *strict-subtuple-trans'''* **by** *blast*

lemma *subtuple-rtranclp-intro*:
assumes ⋀*xm ym*. *R* (*f xm*) (*f ym*) ⟹ *subtuple R xm ym*
and *bij-on-trancl R f*
and *R*** (*f xm*) (*f ym*)
shows *subtuple R*** *xm ym*
proof −
have (λ*xm ym*. *R* (*f xm*) (*f ym*))** *xm ym*
apply (*insert assms*(*2*) *assms*(*3*))
by (*rule reflect-rtranclp*; *auto*)
with *assms*(*1*) **have** (*subtuple R*)** *xm ym*
by (*metis* (*mono-tags*, *lifting*) *mono-rtranclp*)
hence *subtuple R*** *xm ym*
by (*rule rtrancl-to-subtuple*)
thus *?thesis* **by** *simp*
qed

lemma *strict-subtuple-rtranclp-intro*:
assumes ⋀*xm ym*. *R* (*f xm*) (*f ym*) ⟹
strict-subtuple (λ*x y*. *R x y* ∨ *x* = *y*) *xm ym*
and *bij-on-trancl R f*
and *acyclicP-on* (*fmran' ym*) *R*
and *R*$^{++}$ (*f xm*) (*f ym*)
shows *strict-subtuple R*** *xm ym*
proof −
have (λ*xm ym*. *R* (*f xm*) (*f ym*))$^{++}$ *xm ym*
apply (*insert assms*(*1*) *assms*(*2*) *assms*(*4*))
by (*rule reflect-tranclp*; *auto*)
hence (*strict-subtuple* (λ*x y*. *R x y* ∨ *x* = *y*))$^{++}$ *xm ym*
by (*rule tranclp-trans-induct*;
auto simp add: *assms*(*1*) *tranclp.r-into-trancl*)
with *assms*(*3*) **have** *strict-subtuple* (λ*x y*. *R x y* ∨ *x* = *y*)** *xm ym*
by (*rule trancl-to-strict-subtuple'*)
thus *?thesis* **by** *simp*
qed

1.4.5 Code Setup

abbreviation *subtuple-fun f xm ym* ≡
fBall (*fmdom ym*) (λ*x*. *rel-option f* (*fmlookup xm x*) (*fmlookup ym x*))

abbreviation *strict-subtuple-fun f xm ym* ≡
subtuple-fun f xm ym ∧ *xm* ≠ *ym*

lemma *subtuple-fun-simp* [*code-abbrev*, *simp*]:
subtuple-fun f xm ym = *subtuple f xm ym*
by (*simp add*: *fmrel-on-fset-alt-def*)

lemma *strict-subtuple-fun-simp* [*code-abbrev*, *simp*]:
strict-subtuple-fun f xm ym = *strict-subtuple f xm ym*
by *simp*

end

1.5 Object Model

theory *Object-Model*
imports *HOL−Library.Extended-Nat Finite-Map-Ext*
begin

The section defines a generic object model.

1.5.1 Type Synonyms

type-synonym *attr* = *String.literal*
type-synonym *assoc* = *String.literal*
type-synonym *role* = *String.literal*
type-synonym *oper* = *String.literal*
type-synonym *param* = *String.literal*
type-synonym *elit* = *String.literal*

datatype *param-dir* = *In* | *Out* | *InOut*

type-synonym *'c assoc-end* = *'c* × *nat* × *enat* × *bool* × *bool*
type-synonym *'t param-spec* = *param* × *'t* × *param-dir*
type-synonym (*'t*, *'e*) *oper-spec* =
oper × *'t* × *'t param-spec list* × *'t* × *bool* × *'e option*

definition *assoc-end-class* :: *'c assoc-end* ⇒ *'c* ≡ *fst*
definition *assoc-end-min* :: *'c assoc-end* ⇒ *nat* ≡ *fst* ∘ *snd*
definition *assoc-end-max* :: *'c assoc-end* ⇒ *enat* ≡ *fst* ∘ *snd* ∘ *snd*
definition *assoc-end-ordered* :: *'c assoc-end* ⇒ *bool* ≡ *fst* ∘ *snd* ∘ *snd* ∘ *snd*
definition *assoc-end-unique* :: *'c assoc-end* ⇒ *bool* ≡ *snd* ∘ *snd* ∘ *snd* ∘ *snd*

definition *oper-name* :: (*'t*, *'e*) *oper-spec* ⇒ *oper* ≡ *fst*
definition *oper-context* :: (*'t*, *'e*) *oper-spec* ⇒ *'t* ≡ *fst* ∘ *snd*
definition *oper-params* :: (*'t*, *'e*) *oper-spec* ⇒ *'t param-spec list* ≡ *fst* ∘ *snd* ∘ *snd*

definition *oper-result* :: (*'t*, *'e*) *oper-spec* ⇒ *'t* ≡ *fst* ∘ *snd* ∘ *snd* ∘ *snd*

definition *oper-static* :: ('t, 'e) *oper-spec* ⇒ *bool* ≡ *fst* ∘ *snd* ∘ *snd* ∘ *snd* ∘ *snd*
definition *oper-body* :: ('t, 'e) *oper-spec* ⇒ 'e *option* ≡ *snd* ∘ *snd* ∘ *snd* ∘ *snd* ∘ *snd*

definition *param-name* ::'t *param-spec* ⇒ *param* ≡ *fst*
definition *param-type* ::'t *param-spec* ⇒ 't ≡ *fst* ∘ *snd*
definition *param-dir* ::'t *param-spec* ⇒ *param-dir* ≡ *snd* ∘ *snd*

definition *oper-in-params op* ≡
filter (λp. *param-dir p* = *In* ∨ *param-dir p* = *InOut*) (*oper-params op*)

definition *oper-out-params op* ≡
filter (λp. *param-dir p* = *Out* ∨ *param-dir p* = *InOut*) (*oper-params op*)

1.5.2 Attributes

inductive *owned-attribute'* **where**
$\mathcal{C}$ |∈| *fmdom attributes* ⟹
fmlookup attributes $\mathcal{C}$ = *Some* $attrs_{\mathcal{C}}$ ⟹
fmlookup $attrs_{\mathcal{C}}$ *attr* = *Some* τ ⟹
owned-attribute' attributes $\mathcal{C}$ *attr* τ

inductive *attribute-not-closest* **where**
owned-attribute' attributes $\mathcal{D}'$ *attr* τ' ⟹
$\mathcal{C} \leq \mathcal{D}'$ ⟹
$\mathcal{D}' < \mathcal{D}$ ⟹
attribute-not-closest attributes $\mathcal{C}$ *attr* $\mathcal{D}$ τ

inductive *closest-attribute* **where**
owned-attribute' attributes $\mathcal{D}$ *attr* τ ⟹
$\mathcal{C} \leq \mathcal{D}$ ⟹
¬ *attribute-not-closest attributes* $\mathcal{C}$ *attr* $\mathcal{D}$ τ ⟹
closest-attribute attributes $\mathcal{C}$ *attr* $\mathcal{D}$ τ

inductive *closest-attribute-not-unique* **where**
closest-attribute attributes $\mathcal{C}$ *attr* $\mathcal{D}'$ τ' ⟹
$\mathcal{D} \neq \mathcal{D}' \vee \tau \neq \tau'$ ⟹
closest-attribute-not-unique attributes $\mathcal{C}$ *attr* $\mathcal{D}$ τ

inductive *unique-closest-attribute* **where**
closest-attribute attributes $\mathcal{C}$ *attr* $\mathcal{D}$ τ ⟹
¬ *closest-attribute-not-unique attributes* $\mathcal{C}$ *attr* $\mathcal{D}$ τ ⟹
unique-closest-attribute attributes $\mathcal{C}$ *attr* $\mathcal{D}$ τ

1.5.3 Association Ends

inductive *role-refer-class* **where**
role |∈| *fmdom ends* ⟹
fmlookup ends role = *Some end* ⟹
assoc-end-class end = $\mathcal{C}$ ⟹

role-refer-class ends $\mathcal{C}$ *role*

inductive *association-ends′* **where**
$\mathcal{C}$ |∈| *classes* $\Longrightarrow$
assoc |∈| *fmdom associations* $\Longrightarrow$
fmlookup associations assoc = *Some ends* $\Longrightarrow$
role-refer-class ends $\mathcal{C}$ *from* $\Longrightarrow$
role |∈| *fmdom ends* $\Longrightarrow$
fmlookup ends role = *Some end* $\Longrightarrow$
role $\neq$ *from* $\Longrightarrow$
association-ends′ classes associations $\mathcal{C}$ *from role end*

inductive *association-ends-not-unique′* **where**
association-ends′ classes associations $\mathcal{C}$ *from role* end_1 $\Longrightarrow$
association-ends′ classes associations $\mathcal{C}$ *from role* end_2 $\Longrightarrow$
$end_1 \neq end_2$ $\Longrightarrow$
association-ends-not-unique′ classes associations

inductive *owned-association-end′* **where**
association-ends′ classes associations $\mathcal{C}$ *from role end* $\Longrightarrow$
owned-association-end′ classes associations $\mathcal{C}$ *None role end*
| *association-ends′ classes associations* $\mathcal{C}$ *from role end* $\Longrightarrow$
owned-association-end′ classes associations $\mathcal{C}$ (*Some from*) *role end*

inductive *association-end-not-closest* **where**
owned-association-end′ classes associations $\mathcal{D}'$ *from role end′* $\Longrightarrow$
$\mathcal{C} \leq \mathcal{D}'$ $\Longrightarrow$
$\mathcal{D}' < \mathcal{D}$ $\Longrightarrow$
association-end-not-closest classes associations $\mathcal{C}$ *from role* $\mathcal{D}$ *end*

inductive *closest-association-end* **where**
owned-association-end′ classes associations $\mathcal{D}$ *from role end* $\Longrightarrow$
$\mathcal{C} \leq \mathcal{D}$ $\Longrightarrow$
$\neg$ *association-end-not-closest classes associations* $\mathcal{C}$ *from role* $\mathcal{D}$ *end* $\Longrightarrow$
closest-association-end classes associations $\mathcal{C}$ *from role* $\mathcal{D}$ *end*

inductive *closest-association-end-not-unique* **where**
closest-association-end classes associations $\mathcal{C}$ *from role* $\mathcal{D}'$ *end′* $\Longrightarrow$
$\mathcal{D} \neq \mathcal{D}' \vee end \neq end'$ $\Longrightarrow$
closest-association-end-not-unique classes associations $\mathcal{C}$ *from role* $\mathcal{D}$ *end*

inductive *unique-closest-association-end* **where**
closest-association-end classes associations $\mathcal{C}$ *from role* $\mathcal{D}$ *end* $\Longrightarrow$
$\neg$ *closest-association-end-not-unique classes associations* $\mathcal{C}$ *from role* $\mathcal{D}$ *end* $\Longrightarrow$
unique-closest-association-end classes associations $\mathcal{C}$ *from role* $\mathcal{D}$ *end*

1.5.4 Association Classes

inductive *referred-by-association-class″* **where**

fmlookup association-classes $\mathcal{A}$ = *Some assoc* $\Longrightarrow$
fmlookup associations assoc = *Some ends* $\Longrightarrow$
role-refer-class ends $\mathcal{C}$ *from* $\Longrightarrow$
referred-by-association-class''' association-classes associations $\mathcal{C}$ *from* $\mathcal{A}$

inductive *referred-by-association-class'* **where**
referred-by-association-class''' association-classes associations $\mathcal{C}$ *from* $\mathcal{A} \Longrightarrow$
referred-by-association-class' association-classes associations $\mathcal{C}$ *None* $\mathcal{A}$
| *referred-by-association-class''' association-classes associations* $\mathcal{C}$ *from* $\mathcal{A} \Longrightarrow$
referred-by-association-class' association-classes associations $\mathcal{C}$ (*Some from*) $\mathcal{A}$

inductive *association-class-not-closest* **where**
referred-by-association-class' association-classes associations $\mathcal{D}'$ *from* $\mathcal{A} \Longrightarrow$
$\mathcal{C} \leq \mathcal{D}' \Longrightarrow$
$\mathcal{D}' < \mathcal{D} \Longrightarrow$
association-class-not-closest association-classes associations $\mathcal{C}$ *from* $\mathcal{A}$ $\mathcal{D}$

inductive *closest-association-class* **where**
referred-by-association-class' association-classes associations $\mathcal{D}$ *from* $\mathcal{A} \Longrightarrow$
$\mathcal{C} \leq \mathcal{D} \Longrightarrow$
$\neg$ *association-class-not-closest association-classes associations* $\mathcal{C}$ *from* $\mathcal{A}$ $\mathcal{D} \Longrightarrow$
closest-association-class association-classes associations $\mathcal{C}$ *from* $\mathcal{A}$ $\mathcal{D}$

inductive *closest-association-class-not-unique* **where**
closest-association-class association-classes associations $\mathcal{C}$ *from* $\mathcal{A}$ $\mathcal{D}' \Longrightarrow$
$\mathcal{D} \neq \mathcal{D}' \Longrightarrow$
closest-association-class-not-unique
association-classes associations $\mathcal{C}$ *from* $\mathcal{A}$ $\mathcal{D}$

inductive *unique-closest-association-class* **where**
closest-association-class association-classes associations $\mathcal{C}$ *from* $\mathcal{A}$ $\mathcal{D} \Longrightarrow$
$\neg$ *closest-association-class-not-unique*
association-classes associations $\mathcal{C}$ *from* $\mathcal{A}$ $\mathcal{D} \Longrightarrow$
unique-closest-association-class association-classes associations $\mathcal{C}$ *from* $\mathcal{A}$ $\mathcal{D}$

1.5.5 Association Class Ends

inductive *association-class-end'* **where**
fmlookup association-classes $\mathcal{A}$ = *Some assoc* $\Longrightarrow$
fmlookup associations assoc = *Some ends* $\Longrightarrow$
fmlookup ends role = *Some end* $\Longrightarrow$
association-class-end' association-classes associations $\mathcal{A}$ *role end*

inductive *association-class-end-not-unique* **where**
association-class-end' association-classes associations $\mathcal{A}$ *role end'* $\Longrightarrow$
end $\neq$ *end'* $\Longrightarrow$
association-class-end-not-unique association-classes associations $\mathcal{A}$ *role end*

inductive *unique-association-class-end* **where**

association-class-end′ association-classes associations $\mathcal{A}$ *role end* $\Longrightarrow$
$\neg$ *association-class-end-not-unique*
association-classes associations $\mathcal{A}$ *role end* $\Longrightarrow$
unique-association-class-end association-classes associations $\mathcal{A}$ *role end*

1.5.6 Operations

inductive *any-operation′* **where**
op $|\in|$ *fset-of-list operations* $\Longrightarrow$
oper-name op $=$ *name* $\Longrightarrow$
$\tau \leq$ *oper-context op* $\Longrightarrow$
list-all2 $(\lambda\sigma$ *param.* $\sigma \leq$ *param-type param*$)$ π $($*oper-in-params op*$)$ $\Longrightarrow$
any-operation′ operations τ *name* π *op*

inductive *operation′* **where**
any-operation′ operations τ *name* π *op* $\Longrightarrow$
$\neg$ *oper-static op* $\Longrightarrow$
operation′ operations τ *name* π *op*

inductive *operation-not-unique* **where**
operation′ operations τ *name* π *oper′* $\Longrightarrow$
oper $\neq$ *oper′* $\Longrightarrow$
operation-not-unique operations τ *name* π *oper*

inductive *unique-operation* **where**
operation′ operations τ *name* π *oper* $\Longrightarrow$
$\neg$ *operation-not-unique operations* τ *name* π *oper* $\Longrightarrow$
unique-operation operations τ *name* π *oper*

inductive *operation-defined′* **where**
unique-operation operations τ *name* π *oper* $\Longrightarrow$
operation-defined′ operations τ *name* π

inductive *static-operation′* **where**
any-operation′ operations τ *name* π *op* $\Longrightarrow$
oper-static op $\Longrightarrow$
static-operation′ operations τ *name* π *op*

inductive *static-operation-not-unique* **where**
static-operation′ operations τ *name* π *oper′* $\Longrightarrow$
oper $\neq$ *oper′* $\Longrightarrow$
static-operation-not-unique operations τ *name* π *oper*

inductive *unique-static-operation* **where**
static-operation′ operations τ *name* π *oper* $\Longrightarrow$
$\neg$ *static-operation-not-unique operations* τ *name* π *oper* $\Longrightarrow$
unique-static-operation operations τ *name* π *oper*

inductive *static-operation-defined′* **where**

unique-static-operation operations τ *name* π *oper* $\Longrightarrow$
static-operation-defined′ operations τ *name* π

1.5.7 Literals

inductive *has-literal′* **where**
fmlookup literals e = *Some lits* $\Longrightarrow$
lit $|\in|$ *lits* $\Longrightarrow$
has-literal′ literals e lit

1.5.8 Definition

locale *object-model* =
fixes *classes* :: *'a* :: *semilattice-sup fset*
and *attributes* :: *'a* $\rightharpoonup_f$ *attr* $\rightharpoonup_f$ *'t* :: *order*
and *associations* :: *assoc* $\rightharpoonup_f$ *role* $\rightharpoonup_f$ *'a assoc-end*
and *association-classes* :: *'a* $\rightharpoonup_f$ *assoc*
and *operations* :: (*'t*, *'e*) *oper-spec list*
and *literals* :: *'n* $\rightharpoonup_f$ *elit fset*
assumes *assoc-end-min-less-eq-max*:
assoc $|\in|$ *fmdom associations* $\Longrightarrow$
fmlookup associations assoc = *Some ends* $\Longrightarrow$
role $|\in|$ *fmdom ends* $\Longrightarrow$
fmlookup ends role = *Some end* $\Longrightarrow$
assoc-end-min end $\leq$ *assoc-end-max end*
assumes *association-ends-unique*:
association-ends′ classes associations $\mathcal{C}$ *from role end*$_1$ $\Longrightarrow$
association-ends′ classes associations $\mathcal{C}$ *from role end*$_2$ $\Longrightarrow$ *end*$_1$ = *end*$_2$
begin

abbreviation *owned-attribute* $\equiv$
owned-attribute′ attributes

abbreviation *attribute* $\equiv$
unique-closest-attribute attributes

abbreviation *association-ends* $\equiv$
association-ends′ classes associations

abbreviation *owned-association-end* $\equiv$
owned-association-end′ classes associations

abbreviation *association-end* $\equiv$
unique-closest-association-end classes associations

abbreviation *referred-by-association-class* $\equiv$
unique-closest-association-class association-classes associations

abbreviation *association-class-end* $\equiv$
unique-association-class-end association-classes associations

abbreviation *operation* ≡
unique-operation operations

abbreviation *operation-defined* ≡
operation-defined′ operations

abbreviation *static-operation* ≡
unique-static-operation operations

abbreviation *static-operation-defined* ≡
static-operation-defined′ operations

abbreviation *has-literal* ≡
has-literal′ literals

end

declare *operation-defined′.simps* [*simp*]
declare *static-operation-defined′.simps* [*simp*]

declare *has-literal′.simps* [*simp*]

1.5.9 Properties

lemma (**in** *object-model*) *attribute-det*:
attribute $\mathcal{C}$ *attr* $\mathcal{D}_1$ τ_1 ⟹
attribute $\mathcal{C}$ *attr* $\mathcal{D}_2$ τ_2 ⟹ $\mathcal{D}_1 = \mathcal{D}_2 \wedge \tau_1 = \tau_2$
by (*meson closest-attribute-not-unique.intros unique-closest-attribute.cases*)

lemma (**in** *object-model*) *attribute-self-or-inherited*:
attribute $\mathcal{C}$ *attr* $\mathcal{D}$ τ ⟹ $\mathcal{C} \leq \mathcal{D}$
by (*meson closest-attribute.cases unique-closest-attribute.cases*)

lemma (**in** *object-model*) *attribute-closest*:
attribute $\mathcal{C}$ *attr* $\mathcal{D}$ τ ⟹
owned-attribute $\mathcal{D}'$ *attr* τ ⟹
$\mathcal{C} \leq \mathcal{D}'$ ⟹ ¬ $\mathcal{D}' < \mathcal{D}$
by (*meson attribute-not-closest.intros closest-attribute.cases*
unique-closest-attribute.cases)

lemma (**in** *object-model*) *association-end-det*:
association-end $\mathcal{C}$ *from role* $\mathcal{D}_1$ end_1 ⟹
association-end $\mathcal{C}$ *from role* $\mathcal{D}_2$ end_2 ⟹ $\mathcal{D}_1 = \mathcal{D}_2 \wedge end_1 = end_2$
by (*meson closest-association-end-not-unique.intros*
unique-closest-association-end.cases)

lemma (**in** *object-model*) *association-end-self-or-inherited*:

association-end $\mathcal{C}$ *from role* $\mathcal{D}$ *end* $\Longrightarrow \mathcal{C} \leq \mathcal{D}$
by (*meson closest-association-end.cases unique-closest-association-end.cases*)

lemma (**in** *object-model*) *association-end-closest*:
association-end $\mathcal{C}$ *from role* $\mathcal{D}$ *end* $\Longrightarrow$
owned-association-end $\mathcal{D}'$ *from role end* $\Longrightarrow$
$\mathcal{C} \leq \mathcal{D}' \Longrightarrow \neg\ \mathcal{D}' < \mathcal{D}$
by (*meson association-end-not-closest.intros closest-association-end.cases unique-closest-association-end.cases*)

lemma (**in** *object-model*) *association-class-end-det*:
association-class-end $\mathcal{A}$ *role* $end_1 \Longrightarrow$
association-class-end $\mathcal{A}$ *role* $end_2 \Longrightarrow end_1 = end_2$
by (*meson association-class-end-not-unique.intros unique-association-class-end.cases*)

lemma (**in** *object-model*) *operation-det*:
operation τ *name* π $oper_1 \Longrightarrow$
operation τ *name* π $oper_2 \Longrightarrow oper_1 = oper_2$
by (*meson operation-not-unique.intros unique-operation.cases*)

lemma (**in** *object-model*) *static-operation-det*:
static-operation τ *name* π $oper_1 \Longrightarrow$
static-operation τ *name* π $oper_2 \Longrightarrow oper_1 = oper_2$
by (*meson static-operation-not-unique.intros unique-static-operation.cases*)

1.5.10 Code Setup

declare *owned-attribute'.intros*[*folded Predicate-Compile.contains-def*, *code-pred-intro*]
code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow bool$) *owned-attribute'*
by (*elim owned-attribute'.cases*) (*simp add*: *Predicate-Compile.contains-def*)

code-pred *unique-closest-attribute* **.**

declare *role-refer-class.intros*[*folded Predicate-Compile.contains-def*, *code-pred-intro*]
code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow o \Rightarrow o \Rightarrow bool$) *role-refer-class*
by (*elim role-refer-class.cases*) (*simp add*: *Predicate-Compile.contains-def*)

declare *association-ends'.intros*[*folded Predicate-Compile.contains-def*, *code-pred-intro*]
code-pred (*modes*:

$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow bool)$ *association-ends'*
by (*auto simp*: *Predicate-Compile.contains-def elim*: *association-ends'.cases*)

code-pred *association-ends-not-unique'* **.**

code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow bool$,
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$i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
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$i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow bool)$ *owned-association-end'* **.**

code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$,

$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow o \Rightarrow bool$) *closest-association-end* .

code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow i \Rightarrow i \Rightarrow bool$) *closest-association-end-not-unique* .

code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow o \Rightarrow bool$) *unique-closest-association-end* .

code-pred *unique-closest-association-class* .

code-pred *association-class-end′* .

code-pred *association-class-end-not-unique* .

code-pred *unique-association-class-end* .

declare *any-operation′.intros*[*folded Predicate-Compile.contains-def*, *code-pred-intro*]
code-pred [*show-modes*] *any-operation′*
by (*elim any-operation′.cases*) (*simp add*: *Predicate-Compile.contains-def*)

code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$) *operation′* .

code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$) *operation-not-unique* .

code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$) *unique-operation* .

code-pred *operation-defined′* .

code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$) *static-operation′* .

code-pred (*modes*:

$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$) *static-operation-not-unique* .

code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow bool$,
$i \Rightarrow i \Rightarrow i \Rightarrow i \Rightarrow o \Rightarrow bool$) *unique-static-operation* .

code-pred *static-operation-defined'* .

declare *has-literal'.intros*[*folded Predicate-Compile.contains-def*, *code-pred-intro*]
code-pred (*modes*:
$i \Rightarrow i \Rightarrow i \Rightarrow bool$, $i \Rightarrow i \Rightarrow o \Rightarrow bool$) *has-literal'*
by (*elim has-literal'.cases*) (*simp add*: *Predicate-Compile.contains-def*)

end

Chapter 2

Basic Types

theory *OCL-Basic-Types*
  imports *Main HOL−Library.FSet HOL−Library.Phantom-Type*
begin

2.1 Definition

Basic types are parameterized over classes.

type-synonym *'a enum* = *('a, String.literal) phantom*
type-synonym *elit* = *String.literal*

datatype *('a :: order) basic-type* =
  OclAny
| *OclVoid*
| *Boolean*
| *Real*
| *Integer*
| *UnlimitedNatural*
| *String*
| *ObjectType 'a (‹⟨-⟩$_\mathcal{T}$› [0] 1000)*
| *Enum 'a enum*

inductive *basic-subtype* (**infix** *‹⊏$_B$› 65*) **where**
  OclVoid $\sqsubset_B$ *Boolean*
| *OclVoid* $\sqsubset_B$ *UnlimitedNatural*
| *OclVoid* $\sqsubset_B$ *String*
| *OclVoid* $\sqsubset_B \langle\mathcal{C}\rangle_\mathcal{T}$
| *OclVoid* $\sqsubset_B$ *Enum* $\mathcal{E}$

| *UnlimitedNatural* $\sqsubset_B$ *Integer*
| *Integer* $\sqsubset_B$ *Real*
| $\mathcal{C} < \mathcal{D} \Longrightarrow \langle\mathcal{C}\rangle_\mathcal{T} \sqsubset_B \langle\mathcal{D}\rangle_\mathcal{T}$

| *Boolean* $\sqsubset_B$ *OclAny*
| *Real* $\sqsubset_B$ *OclAny*

| *String* $\sqsubset_B$ *OclAny*
| $\langle \mathcal{C} \rangle_\mathcal{T} \sqsubset_B$ *OclAny*
| *Enum* $\mathcal{E} \sqsubset_B$ *OclAny*

declare *basic-subtype.intros* [*intro*!]

inductive-cases *basic-subtype-x-OclAny* [*elim*!]: $\tau \sqsubset_B$ *OclAny*
inductive-cases *basic-subtype-x-OclVoid* [*elim*!]: $\tau \sqsubset_B$ *OclVoid*
inductive-cases *basic-subtype-x-Boolean* [*elim*!]: $\tau \sqsubset_B$ *Boolean*
inductive-cases *basic-subtype-x-Real* [*elim*!]: $\tau \sqsubset_B$ *Real*
inductive-cases *basic-subtype-x-Integer* [*elim*!]: $\tau \sqsubset_B$ *Integer*
inductive-cases *basic-subtype-x-UnlimitedNatural* [*elim*!]: $\tau \sqsubset_B$ *UnlimitedNatural*

inductive-cases *basic-subtype-x-String* [*elim*!]: $\tau \sqsubset_B$ *String*
inductive-cases *basic-subtype-x-ObjectType* [*elim*!]: $\tau \sqsubset_B \langle \mathcal{C} \rangle_\mathcal{T}$
inductive-cases *basic-subtype-x-Enum* [*elim*!]: $\tau \sqsubset_B$ *Enum* $\mathcal{E}$

inductive-cases *basic-subtype-OclAny-x* [*elim*!]: *OclAny* $\sqsubset_B \sigma$
inductive-cases *basic-subtype-ObjectType-x* [*elim*!]: $\langle \mathcal{C} \rangle_\mathcal{T} \sqsubset_B \sigma$

lemma *basic-subtype-asym*:
$\tau \sqsubset_B \sigma \Longrightarrow \sigma \sqsubset_B \tau \Longrightarrow$ *False*
by (*induct rule*: *basic-subtype.induct*, *auto*)

2.2 Partial Order of Basic Types

instantiation *basic-type* :: (*order*) *order*
begin

definition $(<) \equiv$ *basic-subtype*$^{++}$
definition $(\leq) \equiv$ *basic-subtype***

2.2.1 Strict Introduction Rules

lemma *type-less-x-OclAny-intro* [*intro*]:
$\tau \neq$ *OclAny* $\Longrightarrow \tau <$ *OclAny*
proof −
have *basic-subtype*$^{++}$ *OclVoid OclAny*
by (*rule-tac ?b*= *Boolean* **in** *tranclp.trancl-into-trancl*; *auto*)
moreover have *basic-subtype*$^{++}$ *Integer OclAny*
by (*rule-tac ?b*= *Real* **in** *tranclp.trancl-into-trancl*; *auto*)
moreover hence *basic-subtype*$^{++}$ *UnlimitedNatural OclAny*
by (*rule-tac ?b*= *Integer* **in** *tranclp-into-tranclp2*; *auto*)
ultimately show $\tau \neq$ *OclAny* $\Longrightarrow \tau <$ *OclAny*
unfolding *less-basic-type-def*
by (*induct* τ, *auto*)
qed

lemma *type-less-OclVoid-x-intro* [*intro*]:

$\tau \neq OclVoid \Longrightarrow OclVoid < \tau$
proof −
have *basic-subtype*$^{++}$ *OclVoid OclAny*
by (*rule-tac ?b*= *Boolean* **in** *tranclp.trancl-into-trancl*; *auto*)
moreover have *basic-subtype*$^{++}$ *OclVoid Integer*
by (*rule-tac ?b*= *UnlimitedNatural* **in** *tranclp.trancl-into-trancl*; *auto*)
moreover hence *basic-subtype*$^{++}$ *OclVoid Real*
by (*rule-tac ?b*= *Integer* **in** *tranclp.trancl-into-trancl*; *auto*)
ultimately show $\tau \neq OclVoid \Longrightarrow OclVoid < \tau$
unfolding *less-basic-type-def*
by (*induct* τ; *auto*)
qed

lemma *type-less-x-Real-intro* [*intro*]:
$\tau = UnlimitedNatural \Longrightarrow \tau < Real$
$\tau = Integer \Longrightarrow \tau < Real$
unfolding *less-basic-type-def*
by (*rule rtranclp-into-tranclp2*, *auto*)

lemma *type-less-x-Integer-intro* [*intro*]:
$\tau = UnlimitedNatural \Longrightarrow \tau < Integer$
unfolding *less-basic-type-def*
by (*rule rtranclp-into-tranclp2*, *auto*)

lemma *type-less-x-ObjectType-intro* [*intro*]:
$\tau = \langle \mathcal{C} \rangle_{\mathcal{T}} \Longrightarrow \mathcal{C} < \mathcal{D} \Longrightarrow \tau < \langle \mathcal{D} \rangle_{\mathcal{T}}$
unfolding *less-basic-type-def*
using *dual-order.order-iff-strict* **by** *blast*

2.2.2 Strict Elimination Rules

lemma *type-less-x-OclAny* [*elim*!]:
$\tau < OclAny \Longrightarrow$
$(\tau = OclVoid \Longrightarrow P) \Longrightarrow$
$(\tau = Boolean \Longrightarrow P) \Longrightarrow$
$(\tau = Integer \Longrightarrow P) \Longrightarrow$
$(\tau = UnlimitedNatural \Longrightarrow P) \Longrightarrow$
$(\tau = Real \Longrightarrow P) \Longrightarrow$
$(\tau = String \Longrightarrow P) \Longrightarrow$
$(\bigwedge \mathcal{E}.\ \tau = Enum\ \mathcal{E} \Longrightarrow P) \Longrightarrow$
$(\bigwedge \mathcal{C}.\ \tau = \langle \mathcal{C} \rangle_{\mathcal{T}} \Longrightarrow P) \Longrightarrow P$
unfolding *less-basic-type-def*
by (*induct rule*: *converse-tranclp-induct*; *auto*)

lemma *type-less-x-OclVoid* [*elim*!]:
$\tau < OclVoid \Longrightarrow P$
unfolding *less-basic-type-def*
by (*induct rule*: *converse-tranclp-induct*; *auto*)

lemma *type-less-x-Boolean* [*elim*!]:
$\tau < \mathit{Boolean} \Longrightarrow$
$(\tau = \mathit{OclVoid} \Longrightarrow P) \Longrightarrow P$
unfolding *less-basic-type-def*
by (*induct rule*: *converse-tranclp-induct*; *auto*)

lemma *type-less-x-Real* [*elim*!]:
$\tau < \mathit{Real} \Longrightarrow$
$(\tau = \mathit{OclVoid} \Longrightarrow P) \Longrightarrow$
$(\tau = \mathit{UnlimitedNatural} \Longrightarrow P) \Longrightarrow$
$(\tau = \mathit{Integer} \Longrightarrow P) \Longrightarrow P$
unfolding *less-basic-type-def*
by (*induct rule*: *converse-tranclp-induct*; *auto*)

lemma *type-less-x-Integer* [*elim*!]:
$\tau < \mathit{Integer} \Longrightarrow$
$(\tau = \mathit{OclVoid} \Longrightarrow P) \Longrightarrow$
$(\tau = \mathit{UnlimitedNatural} \Longrightarrow P) \Longrightarrow P$
unfolding *less-basic-type-def*
by (*induct rule*: *converse-tranclp-induct*; *auto*)

lemma *type-less-x-UnlimitedNatural* [*elim*!]:
$\tau < \mathit{UnlimitedNatural} \Longrightarrow$
$(\tau = \mathit{OclVoid} \Longrightarrow P) \Longrightarrow P$
unfolding *less-basic-type-def*
by (*induct rule*: *converse-tranclp-induct*; *auto*)

lemma *type-less-x-String* [*elim*!]:
$\tau < \mathit{String} \Longrightarrow$
$(\tau = \mathit{OclVoid} \Longrightarrow P) \Longrightarrow P$
unfolding *less-basic-type-def*
by (*induct rule*: *converse-tranclp-induct*; *auto*)

lemma *type-less-x-ObjectType* [*elim*!]:
$\tau < \langle \mathcal{D} \rangle_{\mathcal{T}} \Longrightarrow$
$(\tau = \mathit{OclVoid} \Longrightarrow P) \Longrightarrow$
$(\bigwedge \mathcal{C}.\ \tau = \langle \mathcal{C} \rangle_{\mathcal{T}} \Longrightarrow \mathcal{C} < \mathcal{D} \Longrightarrow P) \Longrightarrow P$
unfolding *less-basic-type-def*
apply (*induct rule*: *converse-tranclp-induct*)
apply *auto*[*1*]
using *less-trans* **by** *auto*

lemma *type-less-x-Enum* [*elim*!]:
$\tau < \mathit{Enum}\ \mathcal{E} \Longrightarrow$
$(\tau = \mathit{OclVoid} \Longrightarrow P) \Longrightarrow P$
unfolding *less-basic-type-def*
by (*induct rule*: *converse-tranclp-induct*; *auto*)

2.2.3 Properties

lemma *basic-subtype-irrefl*:
$\tau < \tau \Longrightarrow False$
for τ :: *'a basic-type*
by (*cases* τ; *auto*)

lemma *tranclp-less-basic-type*:
$(\tau, \sigma) \in \{(\tau, \sigma).\ \tau \sqsubset_B \sigma\}^+ \longleftrightarrow \tau < \sigma$
by (*simp add*: *tranclp-unfold less-basic-type-def*)

lemma *basic-subtype-acyclic*:
acyclicP basic-subtype
apply (*rule acyclicI*)
using *OCL-Basic-Types.basic-subtype-irrefl*
OCL-Basic-Types.tranclp-less-basic-type **by** *auto*

lemma *less-le-not-le-basic-type*:
$\tau < \sigma \longleftrightarrow \tau \leq \sigma \wedge \neg\ \sigma \leq \tau$
for τ σ :: *'a basic-type*
unfolding *less-basic-type-def less-eq-basic-type-def*
apply (*rule iffI*; *auto*)
apply (*metis* (*mono-tags*) *basic-subtype-irrefl*
less-basic-type-def tranclp-rtranclp-tranclp)
by (*drule rtranclpD*; *auto*)

lemma *antisym-basic-type*:
$\tau \leq \sigma \Longrightarrow \sigma \leq \tau \Longrightarrow \tau = \sigma$
for τ σ :: *'a basic-type*
unfolding *less-eq-basic-type-def less-basic-type-def*
by (*metis* (*mono-tags*, *lifting*) *less-eq-basic-type-def*
less-le-not-le-basic-type less-basic-type-def rtranclpD)

lemma *order-refl-basic-type* [*iff*]:
$\tau \leq \tau$
for τ :: *'a basic-type*
by (*simp add*: *less-eq-basic-type-def*)

instance
by *standard* (*auto simp add*: *less-eq-basic-type-def*
less-le-not-le-basic-type antisym-basic-type)

end

2.2.4 Non-Strict Introduction Rules

lemma *type-less-eq-x-OclAny-intro* [*intro*]:
$\tau \leq OclAny$
using *order.order-iff-strict* **by** *auto*

lemma *type-less-eq-OclVoid-x-intro* [*intro*]:
$OclVoid \leq \tau$
using *order.order-iff-strict* **by** *auto*

lemma *type-less-eq-x-Real-intro* [*intro*]:
$\tau = UnlimitedNatural \Longrightarrow \tau \leq Real$
$\tau = Integer \Longrightarrow \tau \leq Real$
using *order.order-iff-strict* **by** *auto*

lemma *type-less-eq-x-Integer-intro* [*intro*]:
$\tau = UnlimitedNatural \Longrightarrow \tau \leq Integer$
using *order.order-iff-strict* **by** *auto*

lemma *type-less-eq-x-ObjectType-intro* [*intro*]:
$\tau = \langle \mathcal{C} \rangle_{\mathcal{T}} \Longrightarrow \mathcal{C} \leq \mathcal{D} \Longrightarrow \tau \leq \langle \mathcal{D} \rangle_{\mathcal{T}}$
using *order.order-iff-strict* **by** *fastforce*

2.2.5 Non-Strict Elimination Rules

lemma *type-less-eq-x-OclAny* [*elim*!]:
$\tau \leq OclAny \Longrightarrow$
$(\tau = OclVoid \Longrightarrow P) \Longrightarrow$
$(\tau = OclAny \Longrightarrow P) \Longrightarrow$
$(\tau = Boolean \Longrightarrow P) \Longrightarrow$
$(\tau = Integer \Longrightarrow P) \Longrightarrow$
$(\tau = UnlimitedNatural \Longrightarrow P) \Longrightarrow$
$(\tau = Real \Longrightarrow P) \Longrightarrow$
$(\tau = String \Longrightarrow P) \Longrightarrow$
$(\bigwedge \mathcal{E}.\ \tau = Enum\ \mathcal{E} \Longrightarrow P) \Longrightarrow$
$(\bigwedge \mathcal{C}.\ \tau = \langle \mathcal{C} \rangle_{\mathcal{T}} \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-OclVoid* [*elim*!]:
$\tau \leq OclVoid \Longrightarrow$
$(\tau = OclVoid \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-Boolean* [*elim*!]:
$\tau \leq Boolean \Longrightarrow$
$(\tau = OclVoid \Longrightarrow P) \Longrightarrow$
$(\tau = Boolean \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-Real* [*elim*!]:
$\tau \leq Real \Longrightarrow$
$(\tau = OclVoid \Longrightarrow P) \Longrightarrow$
$(\tau = UnlimitedNatural \Longrightarrow P) \Longrightarrow$
$(\tau = Integer \Longrightarrow P) \Longrightarrow$
$(\tau = Real \Longrightarrow P) \Longrightarrow P$

by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-Integer* [*elim*!]:
$\tau \leq Integer \Longrightarrow$
$(\tau = OclVoid \Longrightarrow P) \Longrightarrow$
$(\tau = UnlimitedNatural \Longrightarrow P) \Longrightarrow$
$(\tau = Integer \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-UnlimitedNatural* [*elim*!]:
$\tau \leq UnlimitedNatural \Longrightarrow$
$(\tau = OclVoid \Longrightarrow P) \Longrightarrow$
$(\tau = UnlimitedNatural \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-String* [*elim*!]:
$\tau \leq String \Longrightarrow$
$(\tau = OclVoid \Longrightarrow P) \Longrightarrow$
$(\tau = String \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-ObjectType* [*elim*!]:
$\tau \leq \langle \mathcal{D} \rangle_{\mathcal{T}} \Longrightarrow$
$(\tau = OclVoid \Longrightarrow P) \Longrightarrow$
$(\bigwedge \mathcal{C}.\ \tau = \langle \mathcal{C} \rangle_{\mathcal{T}} \Longrightarrow \mathcal{C} \leq \mathcal{D} \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-Enum* [*elim*!]:
$\tau \leq Enum\ \mathcal{E} \Longrightarrow$
$(\tau = OclVoid \Longrightarrow P) \Longrightarrow$
$(\tau = Enum\ \mathcal{E} \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

2.2.6 Simplification Rules

lemma *basic-type-less-left-simps* [*simp*]:
$OclAny < \sigma = False$
$OclVoid < \sigma = (\sigma \neq OclVoid)$
$Boolean < \sigma = (\sigma = OclAny)$
$Real < \sigma = (\sigma = OclAny)$
$Integer < \sigma = (\sigma = OclAny \vee \sigma = Real)$
$UnlimitedNatural < \sigma = (\sigma = OclAny \vee \sigma = Real \vee \sigma = Integer)$
$String < \sigma = (\sigma = OclAny)$
$ObjectType\ \mathcal{C} < \sigma = (\exists \mathcal{D}.\ \sigma = OclAny \vee \sigma = ObjectType\ \mathcal{D} \wedge \mathcal{C} < \mathcal{D})$
$Enum\ \mathcal{E} < \sigma = (\sigma = OclAny)$
by (*induct* σ, *auto*)

lemma *basic-type-less-right-simps* [*simp*]:
$\tau < OclAny = (\tau \neq OclAny)$

$\tau < OclVoid = False$
$\tau < Boolean = (\tau = OclVoid)$
$\tau < Real = (\tau = Integer \vee \tau = UnlimitedNatural \vee \tau = OclVoid)$
$\tau < Integer = (\tau = UnlimitedNatural \vee \tau = OclVoid)$
$\tau < UnlimitedNatural = (\tau = OclVoid)$
$\tau < String = (\tau = OclVoid)$
$\tau < ObjectType\ \mathcal{D} = (\exists \mathcal{C}.\ \tau = ObjectType\ \mathcal{C} \wedge \mathcal{C} < \mathcal{D} \vee \tau = OclVoid)$
$\tau < Enum\ \mathcal{E} = (\tau = OclVoid)$
by *auto*

2.3 Upper Semilattice of Basic Types

notation *sup* (**infixl** ‹⊔› *65*)

instantiation *basic-type* :: (*semilattice-sup*) *semilattice-sup*
begin

fun *sup-basic-type* **where**
$\langle \mathcal{C} \rangle_{\mathcal{T}} \sqcup \sigma = (case\ \sigma\ of\ OclVoid \Rightarrow \langle \mathcal{C} \rangle_{\mathcal{T}} \mid \langle \mathcal{D} \rangle_{\mathcal{T}} \Rightarrow \langle \mathcal{C} \sqcup \mathcal{D} \rangle_{\mathcal{T}} \mid \text{-} \Rightarrow OclAny)$
| $\tau \sqcup \sigma = (if\ \tau \leq \sigma\ then\ \sigma\ else\ (if\ \sigma \leq \tau\ then\ \tau\ else\ OclAny))$

lemma *sup-ge1-ObjectType*:
$\langle \mathcal{C} \rangle_{\mathcal{T}} \leq \langle \mathcal{C} \rangle_{\mathcal{T}} \sqcup \sigma$
apply (*induct* σ; *simp add*: *basic-subtype.simps*
less-eq-basic-type-def r-into-rtranclp)
by (*metis Nitpick.rtranclp-unfold basic-subtype.intros*(*8*)
le-imp-less-or-eq r-into-rtranclp sup-ge1)

lemma *sup-ge1-basic-type*:
$\tau \leq \tau \sqcup \sigma$
for $\tau\ \sigma$ :: *'a basic-type*
apply (*induct* τ, *auto*)
using *sup-ge1-ObjectType* **by** *auto*

lemma *sup-commut-basic-type*:
$\tau \sqcup \sigma = \sigma \sqcup \tau$
for $\tau\ \sigma$ :: *'a basic-type*
by (*induct* τ; *induct* σ; *auto simp add*: *sup.commute*)

lemma *sup-least-basic-type*:
$\tau \leq \varrho \Longrightarrow \sigma \leq \varrho \Longrightarrow \tau \sqcup \sigma \leq \varrho$
for $\tau\ \sigma\ \varrho$ :: *'a basic-type*
by (*induct* ϱ; *auto*)

instance
by *standard* (*auto simp add*: *sup-ge1-basic-type*
sup-commut-basic-type sup-least-basic-type)

end

2.4 Code Setup

code-pred *basic-subtype* **.**

fun *basic-subtype-fun* :: *'a::order basic-type* ⇒ *'a basic-type* ⇒ *bool* **where**
  basic-subtype-fun OclAny σ = False
| *basic-subtype-fun OclVoid σ = (σ ≠ OclVoid)*
| *basic-subtype-fun Boolean σ = (σ = OclAny)*
| *basic-subtype-fun Real σ = (σ = OclAny)*
| *basic-subtype-fun Integer σ = (σ = Real ∨ σ = OclAny)*
| *basic-subtype-fun UnlimitedNatural σ = (σ = Integer ∨ σ = Real ∨ σ = OclAny)*

| *basic-subtype-fun String σ = (σ = OclAny)*
| *basic-subtype-fun ⟨C⟩T σ = (case σ*
  of ⟨D⟩T ⇒ C < D
  | OclAny ⇒ True
  | - ⇒ False)
| *basic-subtype-fun (Enum -) σ = (σ = OclAny)*

lemma *less-basic-type-code* [*code*]:
  (<) = basic-subtype-fun
proof (*intro ext iffI*)
 fix *τ σ* :: *'a basic-type*
 show *τ < σ ⟹ basic-subtype-fun τ σ*
  apply (*cases σ*; *auto*)
  using *basic-subtype-fun.elims*(*3*) **by** *fastforce*
 show *basic-subtype-fun τ σ ⟹ τ < σ*
  apply (*erule basic-subtype-fun.elims*, *auto*)
  by (*cases σ*, *auto*)
qed

lemma *less-eq-basic-type-code* [*code*]:
  (≤) = (λx y. basic-subtype-fun x y ∨ x = y)
 unfolding *dual-order.order-iff-strict less-basic-type-code*
 by *auto*

end

Chapter 3

Types

theory *OCL-Types*
imports *OCL-Basic-Types Errorable Tuple*
begin

3.1 Definition

Types are parameterized over classes.

type-synonym *telem* = *String.literal*

datatype (*plugins del*: *size*) *'a type* =
OclSuper
| *Required 'a basic-type* (‹-[*1*]› [*1000*] *1000*)
| *Optional 'a basic-type* (‹-[*?*]› [*1000*] *1000*)
| *Collection 'a type*
| *Set 'a type*
| *OrderedSet 'a type*
| *Bag 'a type*
| *Sequence 'a type*
| *Tuple telem* $\rightharpoonup_f$ *'a type*

We define the *OclInvalid* type separately, because we do not need types like *Set*(*OclInvalid*) in the theory. The *OclVoid*[*1*] type is not equal to *OclInvalid*. For example, *Set*(*OclVoid*[*1*]) could theoretically be a valid type of the following expression:

```
Set{null}->excluding(null)
```

definition *OclInvalid* :: *'a type*$_\perp$ $\equiv \perp$

instantiation *type* :: (*type*) *size*
begin

primrec *size-type* :: *'a type* ⇒ *nat* **where**
size-type OclSuper = *0*
| *size-type* (*Required* τ) = *0*

| *size-type* (*Optional* τ) = *0*
| *size-type* (*Collection* τ) = *Suc* (*size-type* τ)
| *size-type* (*Set* τ) = *Suc* (*size-type* τ)
| *size-type* (*OrderedSet* τ) = *Suc* (*size-type* τ)
| *size-type* (*Bag* τ) = *Suc* (*size-type* τ)
| *size-type* (*Sequence* τ) = *Suc* (*size-type* τ)
| *size-type* (*Tuple* π) = *Suc* (*ffold tcf 0* (*fset-of-fmap* (*fmmap size-type* π)))

instance ..

end

inductive *subtype* :: *'a*::*order type* ⇒ *'a type* ⇒ *bool* (**infix** ‹⊏› *65*) **where**
  $\tau \sqsubset_B \sigma \Longrightarrow \tau[1] \sqsubset \sigma[1]$
| $\tau \sqsubset_B \sigma \Longrightarrow \tau[?] \sqsubset \sigma[?]$
| $\tau[1] \sqsubset \tau[?]$
| *OclAny*$[?] \sqsubset$ *OclSuper*

| $\tau \sqsubset \sigma \Longrightarrow$ *Collection* $\tau \sqsubset$ *Collection* σ
| $\tau \sqsubset \sigma \Longrightarrow$ *Set* $\tau \sqsubset$ *Set* σ
| $\tau \sqsubset \sigma \Longrightarrow$ *OrderedSet* $\tau \sqsubset$ *OrderedSet* σ
| $\tau \sqsubset \sigma \Longrightarrow$ *Bag* $\tau \sqsubset$ *Bag* σ
| $\tau \sqsubset \sigma \Longrightarrow$ *Sequence* $\tau \sqsubset$ *Sequence* σ
| *Set* $\tau \sqsubset$ *Collection* τ
| *OrderedSet* $\tau \sqsubset$ *Collection* τ
| *Bag* $\tau \sqsubset$ *Collection* τ
| *Sequence* $\tau \sqsubset$ *Collection* τ
| *Collection OclSuper* $\sqsubset$ *OclSuper*

| *strict-subtuple* $(\lambda\tau\ \sigma.\ \tau \sqsubset \sigma \vee \tau = \sigma)\ \pi\ \xi \Longrightarrow$
  Tuple $\pi \sqsubset$ *Tuple* ξ
| *Tuple* $\pi \sqsubset$ *OclSuper*

declare *subtype.intros* [*intro*!]

inductive-cases *subtype-x-OclSuper* [*elim*!]: $\tau \sqsubset$ *OclSuper*
inductive-cases *subtype-x-Required* [*elim*!]: $\tau \sqsubset \sigma[1]$
inductive-cases *subtype-x-Optional* [*elim*!]: $\tau \sqsubset \sigma[?]$
inductive-cases *subtype-x-Collection* [*elim*!]: $\tau \sqsubset$ *Collection* σ
inductive-cases *subtype-x-Set* [*elim*!]: $\tau \sqsubset$ *Set* σ
inductive-cases *subtype-x-OrderedSet* [*elim*!]: $\tau \sqsubset$ *OrderedSet* σ
inductive-cases *subtype-x-Bag* [*elim*!]: $\tau \sqsubset$ *Bag* σ
inductive-cases *subtype-x-Sequence* [*elim*!]: $\tau \sqsubset$ *Sequence* σ
inductive-cases *subtype-x-Tuple* [*elim*!]: $\tau \sqsubset$ *Tuple* π

inductive-cases *subtype-OclSuper-x* [*elim*!]: *OclSuper* $\sqsubset \sigma$
inductive-cases *subtype-Collection-x* [*elim*!]: *Collection* $\tau \sqsubset \sigma$

lemma *subtype-asym*:

$\tau \sqsubset \sigma \Longrightarrow \sigma \sqsubset \tau \Longrightarrow \mathit{False}$
apply (*induct rule*: *subtype.induct*)
using *basic-subtype-asym* **apply** *auto*
using *subtuple-antisym* **by** *fastforce*

3.2 Constructors Bijectivity on Transitive Closures

lemma *Required-bij-on-trancl* [*simp*]:
bij-on-trancl subtype Required
by (*auto simp add*: *inj-def*)

lemma *not-subtype-Optional-Required*:
$\mathit{subtype}^{++}\ \tau[?]\ \sigma \Longrightarrow \sigma = \varrho[1] \Longrightarrow P$
by (*induct arbitrary*: ϱ *rule*: *tranclp-induct*; *auto*)

lemma *Optional-bij-on-trancl* [*simp*]:
bij-on-trancl subtype Optional
apply (*auto simp add*: *inj-def*)
using *not-subtype-Optional-Required* **by** *blast*

lemma *subtype-tranclp-Collection-x*:
$\mathit{subtype}^{++}\ (\mathit{Collection}\ \tau)\ \sigma \Longrightarrow$
$(\bigwedge \varrho.\ \sigma = \mathit{Collection}\ \varrho \Longrightarrow \mathit{subtype}^{++}\ \tau\ \varrho \Longrightarrow P) \Longrightarrow$
$(\sigma = \mathit{OclSuper} \Longrightarrow P) \Longrightarrow P$
apply (*induct rule*: *tranclp-induct*, *auto*)
by (*metis subtype-Collection-x subtype-OclSuper-x tranclp.trancl-into-trancl*)

lemma *Collection-bij-on-trancl* [*simp*]:
bij-on-trancl subtype Collection
apply (*auto simp add*: *inj-def*)
using *subtype-tranclp-Collection-x* **by** *auto*

lemma *Set-bij-on-trancl* [*simp*]:
bij-on-trancl subtype Set
by (*auto simp add*: *inj-def*)

lemma *OrderedSet-bij-on-trancl* [*simp*]:
bij-on-trancl subtype OrderedSet
by (*auto simp add*: *inj-def*)

lemma *Bag-bij-on-trancl* [*simp*]:
bij-on-trancl subtype Bag
by (*auto simp add*: *inj-def*)

lemma *Sequence-bij-on-trancl* [*simp*]:
bij-on-trancl subtype Sequence
by (*auto simp add*: *inj-def*)

lemma *Tuple-bij-on-trancl* [*simp*]:

bij-on-trancl subtype Tuple
by (*auto simp add*: *inj-def*)

3.3 Partial Order of Types

instantiation *type* :: (*order*) *order*
begin

definition $(<) \equiv subtype^{++}$
definition $(\leq) \equiv subtype^{**}$

3.3.1 Strict Introduction Rules

lemma *type-less-x-Required-intro* [*intro*]:
$\tau = \varrho[1] \Longrightarrow \varrho < \sigma \Longrightarrow \tau < \sigma[1]$
unfolding *less-type-def less-basic-type-def*
by *simp* (*rule preserve-tranclp*; *auto*)

lemma *type-less-x-Optional-intro* [*intro*]:
$\tau = \varrho[1] \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau < \sigma[?]$
$\tau = \varrho[?] \Longrightarrow \varrho < \sigma \Longrightarrow \tau < \sigma[?]$
unfolding *less-type-def less-basic-type-def less-eq-basic-type-def*
apply *simp-all*
apply (*rule preserve-rtranclp''*; *auto*)
by (*rule preserve-tranclp*; *auto*)

lemma *type-less-x-Collection-intro* [*intro*]:
$\tau = Collection\ \varrho \Longrightarrow \varrho < \sigma \Longrightarrow \tau < Collection\ \sigma$
$\tau = Set\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau < Collection\ \sigma$
$\tau = OrderedSet\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau < Collection\ \sigma$
$\tau = Bag\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau < Collection\ \sigma$
$\tau = Sequence\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau < Collection\ \sigma$
unfolding *less-type-def less-eq-type-def*
apply *simp-all*
apply (*rule-tac ?f*= *Collection* **in** *preserve-tranclp*; *auto*)
apply (*rule preserve-rtranclp''*; *auto*)
apply (*rule preserve-rtranclp''*; *auto*)
apply (*rule preserve-rtranclp''*; *auto*)
by (*rule preserve-rtranclp''*; *auto*)

lemma *type-less-x-Set-intro* [*intro*]:
$\tau = Set\ \varrho \Longrightarrow \varrho < \sigma \Longrightarrow \tau < Set\ \sigma$
unfolding *less-type-def*
by *simp* (*rule preserve-tranclp*; *auto*)

lemma *type-less-x-OrderedSet-intro* [*intro*]:
$\tau = OrderedSet\ \varrho \Longrightarrow \varrho < \sigma \Longrightarrow \tau < OrderedSet\ \sigma$
unfolding *less-type-def*
by *simp* (*rule preserve-tranclp*; *auto*)

lemma *type-less-x-Bag-intro* [*intro*]:
$\tau = Bag\ \varrho \Longrightarrow \varrho < \sigma \Longrightarrow \tau < Bag\ \sigma$
unfolding *less-type-def*
by *simp* (*rule preserve-tranclp*; *auto*)

lemma *type-less-x-Sequence-intro* [*intro*]:
$\tau = Sequence\ \varrho \Longrightarrow \varrho < \sigma \Longrightarrow \tau < Sequence\ \sigma$
unfolding *less-type-def*
by *simp* (*rule preserve-tranclp*; *auto*)

lemma *fun-or-eq-refl* [*intro*]:
reflp $(\lambda x\ y.\ f\ x\ y \vee x = y)$
by (*simp add*: *reflpI*)

lemma *type-less-x-Tuple-intro* [*intro*]:
assumes $\tau = Tuple\ \pi$
and *strict-subtuple* $(\leq)\ \pi\ \xi$
shows $\tau < Tuple\ \xi$
proof −
have *subtuple* $(\lambda\tau\ \sigma.\ \tau \sqsubset \sigma \vee \tau = \sigma)^{**}\ \pi\ \xi$
using *assms*(*2*) *less-eq-type-def* **by** *auto*
hence $(subtuple\ (\lambda\tau\ \sigma.\ \tau \sqsubset \sigma \vee \tau = \sigma))^{++}\ \pi\ \xi$
by *simp* (*rule subtuple-to-trancl*; *auto*)
hence $(strict\text{-}subtuple\ (\lambda\tau\ \sigma.\ \tau \sqsubset \sigma \vee \tau = \sigma))^{**}\ \pi\ \xi$
by (*simp add*: *tranclp-into-rtranclp*)
hence $(strict\text{-}subtuple\ (\lambda\tau\ \sigma.\ \tau \sqsubset \sigma \vee \tau = \sigma))^{++}\ \pi\ \xi$
by (*meson assms*(*2*) *rtranclpD*)
thus *?thesis*
unfolding *less-type-def*
using *assms*(*1*) **apply** *simp*
by (*rule preserve-tranclp*; *auto*)
qed

lemma *type-less-x-OclSuper-intro* [*intro*]:
$\tau \neq OclSuper \Longrightarrow \tau < OclSuper$
unfolding *less-type-def*
proof (*induct* τ)
case *OclSuper* **thus** *?case* **by** *auto*
next
case (*Required* τ)
have $subtype^{**}\ \tau[1]\ OclAny[1]$
apply (*rule-tac ?f*= *Required* **in** *preserve-rtranclp*[*of basic-subtype*], *auto*)
by (*metis less-eq-basic-type-def type-less-eq-x-OclAny-intro*)
also have $subtype^{++}\ OclAny[1]\ OclAny[?]$ **by** *auto*
also have $subtype^{++}\ OclAny[?]\ OclSuper$ **by** *auto*
finally show *?case* **by** *auto*
next
case (*Optional* τ)

```
  have subtype** τ[?] OclAny[?]
    apply (rule-tac ?f= Optional in preserve-rtranclp[of basic-subtype], auto)
    by (metis less-eq-basic-type-def type-less-eq-x-OclAny-intro)
  also have subtype++ OclAny[?] OclSuper by auto
  finally show ?case by auto
next
  case (Collection τ)
  have subtype** (Collection τ) (Collection OclSuper)
    apply (rule-tac ?f= Collection in preserve-rtranclp[of subtype], auto)
    using Collection.hyps by force
  also have subtype++ (Collection OclSuper) OclSuper by auto
  finally show ?case by auto
next
  case (Set τ)
  have subtype++ (Set τ) (Collection τ) by auto
  also have subtype** (Collection τ) (Collection OclSuper)
    apply (rule-tac ?f= Collection in preserve-rtranclp[of subtype], auto)
    using Set.hyps by force
  also have subtype** (Collection OclSuper) OclSuper by auto
  finally show ?case by auto
next
  case (OrderedSet τ)
  have subtype++ (OrderedSet τ) (Collection τ) by auto
  also have subtype** (Collection τ) (Collection OclSuper)
    apply (rule-tac ?f= Collection in preserve-rtranclp[of subtype], auto)
    using OrderedSet.hyps by force
  also have subtype** (Collection OclSuper) OclSuper by auto
  finally show ?case by auto
next
  case (Bag τ)
  have subtype++ (Bag τ) (Collection τ) by auto
  also have subtype** (Collection τ) (Collection OclSuper)
    apply (rule-tac ?f= Collection in preserve-rtranclp[of subtype], auto)
    using Bag.hyps by force
  also have subtype** (Collection OclSuper) OclSuper by auto
  finally show ?case by auto
next
  case (Sequence τ)
  have subtype++ (Sequence τ) (Collection τ) by auto
  also have subtype** (Collection τ) (Collection OclSuper)
    apply (rule-tac ?f= Collection in preserve-rtranclp[of subtype], auto)
    using Sequence.hyps by force
  also have subtype** (Collection OclSuper) OclSuper by auto
  finally show ?case by auto
next
  case (Tuple x) thus ?case by auto
qed
```

3.3.2 Strict Elimination Rules

lemma *type-less-x-Required* [*elim*!]:
assumes $\tau < \sigma[1]$
and $\bigwedge\varrho.\ \tau = \varrho[1] \Longrightarrow \varrho < \sigma \Longrightarrow P$
shows P
proof −
from *assms*(*1*) **obtain** ϱ **where** $\tau = \varrho[1]$
unfolding *less-type-def*
by (*induct rule*: *converse-tranclp-induct*; *auto*)
moreover have $\bigwedge\tau\ \sigma.\ \tau[1] < \sigma[1] \Longrightarrow \tau < \sigma$
unfolding *less-type-def less-basic-type-def*
by (*rule reflect-tranclp*; *auto*)
ultimately show *?thesis*
using *assms* **by** *auto*
qed

lemma *type-less-x-Optional* [*elim*!]:
$\tau < \sigma[?] \Longrightarrow$
$(\bigwedge\varrho.\ \tau = \varrho[1] \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow$
$(\bigwedge\varrho.\ \tau = \varrho[?] \Longrightarrow \varrho < \sigma \Longrightarrow P) \Longrightarrow P$
unfolding *less-type-def*
apply (*induct rule*: *converse-tranclp-induct*)
apply (*metis subtype-x-Optional eq-refl less-basic-type-def tranclp.r-into-trancl*)
apply (*erule subtype.cases*; *auto*)
apply (*simp add*: *converse-rtranclp-into-rtranclp less-eq-basic-type-def*)
by (*simp add*: *less-basic-type-def tranclp-into-tranclp2*)

lemma *type-less-x-Collection* [*elim*!]:
$\tau < Collection\ \sigma \Longrightarrow$
$(\bigwedge\varrho.\ \tau = Collection\ \varrho \Longrightarrow \varrho < \sigma \Longrightarrow P) \Longrightarrow$
$(\bigwedge\varrho.\ \tau = Set\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow$
$(\bigwedge\varrho.\ \tau = OrderedSet\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow$
$(\bigwedge\varrho.\ \tau = Bag\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow$
$(\bigwedge\varrho.\ \tau = Sequence\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow P$
unfolding *less-type-def*
apply (*induct rule*: *converse-tranclp-induct*)
apply (*metis* (*mono-tags*) *Nitpick.rtranclp-unfold*
subtype-x-Collection less-eq-type-def tranclp.r-into-trancl)
by (*erule subtype.cases*;
auto simp add: *converse-rtranclp-into-rtranclp less-eq-type-def*
tranclp-into-tranclp2 tranclp-into-rtranclp)

lemma *type-less-x-Set* [*elim*!]:
assumes $\tau < Set\ \sigma$
and $\bigwedge\varrho.\ \tau = Set\ \varrho \Longrightarrow \varrho < \sigma \Longrightarrow P$
shows P
proof −
from *assms*(*1*) **obtain** ϱ **where** $\tau = Set\ \varrho$
unfolding *less-type-def*

```
    by (induct rule: converse-tranclp-induct; auto)
  moreover have ⋀τ σ. Set τ < Set σ ⟹ τ < σ
    unfolding less-type-def
    by (rule reflect-tranclp; auto)
  ultimately show ?thesis
    using assms by auto
qed

lemma type-less-x-OrderedSet [elim!]:
  assumes τ < OrderedSet σ
      and ⋀ϱ. τ = OrderedSet ϱ ⟹ ϱ < σ ⟹ P
    shows P
proof −
  from assms(1) obtain ϱ where τ = OrderedSet ϱ
    unfolding less-type-def
    by (induct rule: converse-tranclp-induct; auto)
  moreover have ⋀τ σ. OrderedSet τ < OrderedSet σ ⟹ τ < σ
    unfolding less-type-def
    by (rule reflect-tranclp; auto)
  ultimately show ?thesis
    using assms by auto
qed

lemma type-less-x-Bag [elim!]:
  assumes τ < Bag σ
      and ⋀ϱ. τ = Bag ϱ ⟹ ϱ < σ ⟹ P
    shows P
proof −
  from assms(1) obtain ϱ where τ = Bag ϱ
    unfolding less-type-def
    by (induct rule: converse-tranclp-induct; auto)
  moreover have ⋀τ σ. Bag τ < Bag σ ⟹ τ < σ
    unfolding less-type-def
    by (rule reflect-tranclp; auto)
  ultimately show ?thesis
    using assms by auto
qed

lemma type-less-x-Sequence [elim!]:
  assumes τ < Sequence σ
      and ⋀ϱ. τ = Sequence ϱ ⟹ ϱ < σ ⟹ P
    shows P
proof −
  from assms(1) obtain ϱ where τ = Sequence ϱ
    unfolding less-type-def
    by (induct rule: converse-tranclp-induct; auto)
  moreover have ⋀τ σ. Sequence τ < Sequence σ ⟹ τ < σ
    unfolding less-type-def
    by (rule reflect-tranclp; auto)
```

ultimately show *?thesis*
using *assms* **by** *auto*
qed

We will be able to remove the acyclicity assumption only after we prove that the subtype relation is acyclic.

lemma *type-less-x-Tuple'*:
assumes $\tau < Tuple\ \xi$
and *acyclicP-on* (*fmran'* ξ) *subtype*
and $\bigwedge\pi.\ \tau = Tuple\ \pi \Longrightarrow$ *strict-subtuple* $(\leq)\ \pi\ \xi \Longrightarrow P$
shows P
proof −
from *assms*(*1*) **obtain** π **where** $\tau = Tuple\ \pi$
unfolding *less-type-def*
by (*induct rule*: *converse-tranclp-induct*; *auto*)
moreover from *assms*(*2*) **have**
$\bigwedge\pi.\ Tuple\ \pi < Tuple\ \xi \Longrightarrow$ *strict-subtuple* $(\leq)\ \pi\ \xi$
unfolding *less-type-def less-eq-type-def*
by (*rule-tac ?f*= *Tuple* **in** *strict-subtuple-rtranclp-intro*; *auto*)
ultimately show *?thesis*
using *assms* **by** *auto*
qed

lemma *type-less-x-OclSuper* [*elim*!]:
$\tau < OclSuper \Longrightarrow (\tau \neq OclSuper \Longrightarrow P) \Longrightarrow P$
unfolding *less-type-def*
by (*drule tranclpD*, *auto*)

3.3.3 Properties

lemma *subtype-irrefl*:
$\tau < \tau \Longrightarrow False$
for τ :: *'a type*
apply (*induct* τ, *auto*)
apply (*erule type-less-x-Tuple'*, *auto*)
unfolding *less-type-def tranclp-unfold*
by *auto*

lemma *subtype-acyclic*:
acyclicP subtype
apply (*rule acyclicI*)
apply (*simp add*: *trancl-def*)
by (*metis* (*mono-tags*) *OCL-Types.less-type-def OCL-Types.subtype-irrefl*)

lemma *less-le-not-le-type*:
$\tau < \sigma \longleftrightarrow \tau \leq \sigma \wedge \neg\ \sigma \leq \tau$
for τ σ :: *'a type*
proof
show $\tau < \sigma \Longrightarrow \tau \leq \sigma \wedge \neg\ \sigma \leq \tau$

apply (*auto simp add*: *less-type-def less-eq-type-def*)
by (*metis* (*mono-tags*) *subtype-irrefl less-type-def tranclp-rtranclp-tranclp*)
show $\tau \leq \sigma \wedge \neg\ \sigma \leq \tau \Longrightarrow \tau < \sigma$
apply (*auto simp add*: *less-type-def less-eq-type-def*)
by (*metis rtranclpD*)
qed

lemma *order-refl-type* [*iff*]:
$\tau \leq \tau$
for τ :: *'a type*
unfolding *less-eq-type-def* **by** *simp*

lemma *order-trans-type*:
$\tau \leq \sigma \Longrightarrow \sigma \leq \varrho \Longrightarrow \tau \leq \varrho$
for $\tau\ \sigma\ \varrho$ :: *'a type*
unfolding *less-eq-type-def* **by** *simp*

lemma *antisym-type*:
$\tau \leq \sigma \Longrightarrow \sigma \leq \tau \Longrightarrow \tau = \sigma$
for $\tau\ \sigma$ :: *'a type*
unfolding *less-eq-type-def less-type-def*
by (*metis* (*mono-tags*, *lifting*) *less-eq-type-def*
less-le-not-le-type less-type-def rtranclpD)

instance
apply *intro-classes*
apply (*simp add*: *less-le-not-le-type*)
apply (*simp*)
using *order-trans-type* **apply** *blast*
by (*simp add*: *antisym-type*)

end

3.3.4 Non-Strict Introduction Rules

lemma *type-less-eq-x-Required-intro* [*intro*]:
$\tau = \varrho[1] \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq \sigma[1]$
unfolding *dual-order.order-iff-strict* **by** *auto*

lemma *type-less-eq-x-Optional-intro* [*intro*]:
$\tau = \varrho[1] \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq \sigma[?]$
$\tau = \varrho[?] \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq \sigma[?]$
unfolding *dual-order.order-iff-strict* **by** *auto*

lemma *type-less-eq-x-Collection-intro* [*intro*]:
$\tau = Collection\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq Collection\ \sigma$
$\tau = Set\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq Collection\ \sigma$
$\tau = OrderedSet\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq Collection\ \sigma$
$\tau = Bag\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq Collection\ \sigma$

$\tau = Sequence\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq Collection\ \sigma$
unfolding *dual-order.order-iff-strict* **by** *auto*

lemma *type-less-eq-x-Set-intro* [*intro*]:
$\tau = Set\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq Set\ \sigma$
unfolding *dual-order.order-iff-strict* **by** *auto*

lemma *type-less-eq-x-OrderedSet-intro* [*intro*]:
$\tau = OrderedSet\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq OrderedSet\ \sigma$
unfolding *dual-order.order-iff-strict* **by** *auto*

lemma *type-less-eq-x-Bag-intro* [*intro*]:
$\tau = Bag\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq Bag\ \sigma$
unfolding *dual-order.order-iff-strict* **by** *auto*

lemma *type-less-eq-x-Sequence-intro* [*intro*]:
$\tau = Sequence\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow \tau \leq Sequence\ \sigma$
unfolding *dual-order.order-iff-strict* **by** *auto*

lemma *type-less-eq-x-Tuple-intro* [*intro*]:
$\tau = Tuple\ \pi \Longrightarrow subtuple\ (\leq)\ \pi\ \xi \Longrightarrow \tau \leq Tuple\ \xi$
using *dual-order.strict-iff-order* **by** *blast*

lemma *type-less-eq-x-OclSuper-intro* [*intro*]:
$\tau \leq OclSuper$
unfolding *dual-order.order-iff-strict* **by** *auto*

3.3.5 Non-Strict Elimination Rules

lemma *type-less-eq-x-Required* [*elim*!]:
$\tau \leq \sigma[1] \Longrightarrow$
$(\bigwedge\varrho.\ \tau = \varrho[1] \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-Optional* [*elim*!]:
$\tau \leq \sigma[?] \Longrightarrow$
$(\bigwedge\varrho.\ \tau = \varrho[1] \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow$
$(\bigwedge\varrho.\ \tau = \varrho[?] \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*, *auto*)

lemma *type-less-eq-x-Collection* [*elim*!]:
$\tau \leq Collection\ \sigma \Longrightarrow$
$(\bigwedge\varrho.\ \tau = Set\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow$
$(\bigwedge\varrho.\ \tau = OrderedSet\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow$
$(\bigwedge\varrho.\ \tau = Bag\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow$
$(\bigwedge\varrho.\ \tau = Sequence\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow$
$(\bigwedge\varrho.\ \tau = Collection\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-Set* [*elim*!]:
$\tau \leq Set\ \sigma \Longrightarrow$
$(\bigwedge \varrho.\ \tau = Set\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-OrderedSet* [*elim*!]:
$\tau \leq OrderedSet\ \sigma \Longrightarrow$
$(\bigwedge \varrho.\ \tau = OrderedSet\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-Bag* [*elim*!]:
$\tau \leq Bag\ \sigma \Longrightarrow$
$(\bigwedge \varrho.\ \tau = Bag\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-eq-x-Sequence* [*elim*!]:
$\tau \leq Sequence\ \sigma \Longrightarrow$
$(\bigwedge \varrho.\ \tau = Sequence\ \varrho \Longrightarrow \varrho \leq \sigma \Longrightarrow P) \Longrightarrow P$
by (*drule le-imp-less-or-eq*; *auto*)

lemma *type-less-x-Tuple* [*elim*!]:
$\tau < Tuple\ \xi \Longrightarrow$
$(\bigwedge \pi.\ \tau = Tuple\ \pi \Longrightarrow strict\text{-}subtuple\ (\leq)\ \pi\ \xi \Longrightarrow P) \Longrightarrow P$
apply (*erule type-less-x-Tuple'*)
by (*meson acyclic-def subtype-acyclic*)

lemma *type-less-eq-x-Tuple* [*elim*!]:
$\tau \leq Tuple\ \xi \Longrightarrow$
$(\bigwedge \pi.\ \tau = Tuple\ \pi \Longrightarrow subtuple\ (\leq)\ \pi\ \xi \Longrightarrow P) \Longrightarrow P$
apply (*drule le-imp-less-or-eq*, *auto*)
by (*simp add*: *fmap.rel-refl fmrel-to-subtuple*)

3.3.6 Simplification Rules

lemma *type-less-left-simps* [*simp*]:
$OclSuper < \sigma = False$
$\varrho[1] < \sigma = (\exists \upsilon.$
$\sigma = OclSuper \vee$
$\sigma = \upsilon[1] \wedge \varrho < \upsilon \vee$
$\sigma = \upsilon[?] \wedge \varrho \leq \upsilon)$
$\varrho[?] < \sigma = (\exists \upsilon.$
$\sigma = OclSuper \vee$
$\sigma = \upsilon[?] \wedge \varrho < \upsilon)$
$Collection\ \tau < \sigma = (\exists \varphi.$
$\sigma = OclSuper \vee$
$\sigma = Collection\ \varphi \wedge \tau < \varphi)$
$Set\ \tau < \sigma = (\exists \varphi.$
$\sigma = OclSuper \vee$
$\sigma = Collection\ \varphi \wedge \tau \leq \varphi \vee$

```
    σ = Set φ ∧ τ < φ)
  OrderedSet τ < σ = (∃φ.
    σ = OclSuper ∨
    σ = Collection φ ∧ τ ≤ φ ∨
    σ = OrderedSet φ ∧ τ < φ)
  Bag τ < σ = (∃φ.
    σ = OclSuper ∨
    σ = Collection φ ∧ τ ≤ φ ∨
    σ = Bag φ ∧ τ < φ)
  Sequence τ < σ = (∃φ.
    σ = OclSuper ∨
    σ = Collection φ ∧ τ ≤ φ ∨
    σ = Sequence φ ∧ τ < φ)
  Tuple π < σ = (∃ξ.
    σ = OclSuper ∨
    σ = Tuple ξ ∧ strict-subtuple (≤) π ξ)
 by (induct σ; auto)+

lemma type-less-right-simps [simp]:
  τ < OclSuper = (τ ≠ OclSuper)
  τ < υ[1] = (∃ϱ. τ = ϱ[1] ∧ ϱ < υ)
  τ < υ[?] = (∃ϱ. τ = ϱ[1] ∧ ϱ ≤ υ ∨ τ = ϱ[?] ∧ ϱ < υ)
  τ < Collection σ = (∃φ.
    τ = Collection φ ∧ φ < σ ∨
    τ = Set φ ∧ φ ≤ σ ∨
    τ = OrderedSet φ ∧ φ ≤ σ ∨
    τ = Bag φ ∧ φ ≤ σ ∨
    τ = Sequence φ ∧ φ ≤ σ)
  τ < Set σ = (∃φ. τ = Set φ ∧ φ < σ)
  τ < OrderedSet σ = (∃φ. τ = OrderedSet φ ∧ φ < σ)
  τ < Bag σ = (∃φ. τ = Bag φ ∧ φ < σ)
  τ < Sequence σ = (∃φ. τ = Sequence φ ∧ φ < σ)
  τ < Tuple ξ = (∃π. τ = Tuple π ∧ strict-subtuple (≤) π ξ)
 by auto
```

3.4 Upper Semilattice of Types

```
instantiation type :: (semilattice-sup) semilattice-sup
begin

fun sup-type where
  OclSuper ⊔ σ = OclSuper
| Required τ ⊔ σ = (case σ
   of ϱ[1] ⇒ (τ ⊔ ϱ)[1]
    | ϱ[?] ⇒ (τ ⊔ ϱ)[?]
    | - ⇒ OclSuper)
| Optional τ ⊔ σ = (case σ
   of ϱ[1] ⇒ (τ ⊔ ϱ)[?]
    | ϱ[?] ⇒ (τ ⊔ ϱ)[?]
```

```
    | - ⇒ OclSuper)
| Collection τ ⊔ σ = (case σ
    of Collection ϱ ⇒ Collection (τ ⊔ ϱ)
    | Set ϱ ⇒ Collection (τ ⊔ ϱ)
    | OrderedSet ϱ ⇒ Collection (τ ⊔ ϱ)
    | Bag ϱ ⇒ Collection (τ ⊔ ϱ)
    | Sequence ϱ ⇒ Collection (τ ⊔ ϱ)
    | - ⇒ OclSuper)
| Set τ ⊔ σ = (case σ
    of Collection ϱ ⇒ Collection (τ ⊔ ϱ)
    | Set ϱ ⇒ Set (τ ⊔ ϱ)
    | OrderedSet ϱ ⇒ Collection (τ ⊔ ϱ)
    | Bag ϱ ⇒ Collection (τ ⊔ ϱ)
    | Sequence ϱ ⇒ Collection (τ ⊔ ϱ)
    | - ⇒ OclSuper)
| OrderedSet τ ⊔ σ = (case σ
    of Collection ϱ ⇒ Collection (τ ⊔ ϱ)
    | Set ϱ ⇒ Collection (τ ⊔ ϱ)
    | OrderedSet ϱ ⇒ OrderedSet (τ ⊔ ϱ)
    | Bag ϱ ⇒ Collection (τ ⊔ ϱ)
    | Sequence ϱ ⇒ Collection (τ ⊔ ϱ)
    | - ⇒ OclSuper)
| Bag τ ⊔ σ = (case σ
    of Collection ϱ ⇒ Collection (τ ⊔ ϱ)
    | Set ϱ ⇒ Collection (τ ⊔ ϱ)
    | OrderedSet ϱ ⇒ Collection (τ ⊔ ϱ)
    | Bag ϱ ⇒ Bag (τ ⊔ ϱ)
    | Sequence ϱ ⇒ Collection (τ ⊔ ϱ)
    | - ⇒ OclSuper)
| Sequence τ ⊔ σ = (case σ
    of Collection ϱ ⇒ Collection (τ ⊔ ϱ)
    | Set ϱ ⇒ Collection (τ ⊔ ϱ)
    | OrderedSet ϱ ⇒ Collection (τ ⊔ ϱ)
    | Bag ϱ ⇒ Collection (τ ⊔ ϱ)
    | Sequence ϱ ⇒ Sequence (τ ⊔ ϱ)
    | - ⇒ OclSuper)
| Tuple π ⊔ σ = (case σ
    of Tuple ξ ⇒ Tuple (fmmerge-fun (⊔) π ξ)
    | - ⇒ OclSuper)

lemma sup-ge1-type:
  τ ≤ τ ⊔ σ
 for τ σ :: 'a type
proof (induct τ arbitrary: σ)
 case OclSuper show ?case by simp
 case (Required τ) show ?case by (induct σ; auto)
 case (Optional τ) show ?case by (induct σ; auto)
 case (Collection τ) thus ?case by (induct σ; auto)
 case (Set τ) thus ?case by (induct σ; auto)
```

case (*OrderedSet* τ) **thus** *?case* **by** (*induct* σ; *auto*)
case (*Bag* τ) **thus** *?case* **by** (*induct* σ; *auto*)
case (*Sequence* τ) **thus** *?case* **by** (*induct* σ; *auto*)
next
case (*Tuple* π)
moreover have *Tuple-less-eq-sup*:
$(\bigwedge\tau\ \sigma.\ \tau \in fmran'\ \pi \Longrightarrow \tau \leq \tau \sqcup \sigma) \Longrightarrow$
Tuple $\pi \leq$ *Tuple* $\pi \sqcup \sigma$
by (*cases* σ, *auto*)
ultimately show *?case* **by** (*cases* σ, *auto*)
qed

lemma *sup-commut-type*:
$\tau \sqcup \sigma = \sigma \sqcup \tau$
for $\tau\ \sigma$ :: *'a type*
proof (*induct* τ *arbitrary*: σ)
case *OclSuper* **show** *?case* **by** (*cases* σ; *simp add*: *less-eq-type-def*)
case (*Required* τ) **show** *?case* **by** (*cases* σ; *simp add*: *sup-commute*)
case (*Optional* τ) **show** *?case* **by** (*cases* σ; *simp add*: *sup-commute*)
case (*Collection* τ) **thus** *?case* **by** (*cases* σ; *simp*)
case (*Set* τ) **thus** *?case* **by** (*cases* σ; *simp*)
case (*OrderedSet* τ) **thus** *?case* **by** (*cases* σ; *simp*)
case (*Bag* τ) **thus** *?case* **by** (*cases* σ; *simp*)
case (*Sequence* τ) **thus** *?case* **by** (*cases* σ; *simp*)
next
case (*Tuple* π) **thus** *?case*
apply (*cases* σ; *simp add*: *less-eq-type-def*)
using *fmmerge-commut* **by** *blast*
qed

lemma *sup-least-type*:
$\tau \leq \varrho \Longrightarrow \sigma \leq \varrho \Longrightarrow \tau \sqcup \sigma \leq \varrho$
for $\tau\ \sigma\ \varrho$ :: *'a type*
proof (*induct* ϱ *arbitrary*: $\tau\ \sigma$)
case *OclSuper* **show** *?case* **using** *eq-refl* **by** *auto*
next
case (*Required x*) **show** *?case*
apply (*insert Required*)
by (*erule type-less-eq-x-Required*; *erule type-less-eq-x-Required*; *auto*)
next
case (*Optional x*) **show** *?case*
apply (*insert Optional*)
by (*erule type-less-eq-x-Optional*; *erule type-less-eq-x-Optional*; *auto*)
next
case (*Collection* ϱ) **show** *?case*
apply (*insert Collection*)
by (*erule type-less-eq-x-Collection*; *erule type-less-eq-x-Collection*; *auto*)
next
case (*Set* ϱ) **show** *?case*

apply (*insert Set*)
by (*erule type-less-eq-x-Set*; *erule type-less-eq-x-Set*; *auto*)
next
case (*OrderedSet* ϱ) **show** *?case*
apply (*insert OrderedSet*)
by (*erule type-less-eq-x-OrderedSet*; *erule type-less-eq-x-OrderedSet*; *auto*)
next
case (*Bag* ϱ) **show** *?case*
apply (*insert Bag*)
by (*erule type-less-eq-x-Bag*; *erule type-less-eq-x-Bag*; *auto*)
next
case (*Sequence* ϱ) **thus** *?case*
apply (*insert Sequence*)
by (*erule type-less-eq-x-Sequence*; *erule type-less-eq-x-Sequence*; *auto*)
next
case (*Tuple* π) **show** *?case*
apply (*insert Tuple*)
apply (*erule type-less-eq-x-Tuple*; *erule type-less-eq-x-Tuple*; *auto*)
by (*rule-tac* $?\pi =$ (*fmmerge* $(\sqcup)$ π' π'') **in** *type-less-eq-x-Tuple-intro*;
simp add: *fmrel-on-fset-fmmerge1*)
qed

instance
apply *intro-classes*
apply (*simp add*: *sup-ge1-type*)
apply (*simp add*: *sup-commut-type sup-ge1-type*)
by (*simp add*: *sup-least-type*)

end

3.5 Helper Relations

abbreviation *between* ($\langle$-/ = -—-$\rangle$ [51, 51, 51] 50) **where**
$x = y{-}z \equiv y \leq x \wedge x \leq z$

inductive *element-type* **where**
element-type (*Collection* τ) τ
| *element-type* (*Set* τ) τ
| *element-type* (*OrderedSet* τ) τ
| *element-type* (*Bag* τ) τ
| *element-type* (*Sequence* τ) τ

lemma *element-type-alt-simps*:
element-type τ σ =
(*Collection* $\sigma = \tau \vee$
Set $\sigma = \tau \vee$
OrderedSet $\sigma = \tau \vee$
Bag $\sigma = \tau \vee$
Sequence $\sigma = \tau$)

by (*auto simp add*: *element-type.simps*)

inductive *update-element-type* **where**
update-element-type (*Collection* -) τ (*Collection* τ)
| *update-element-type* (*Set* -) τ (*Set* τ)
| *update-element-type* (*OrderedSet* -) τ (*OrderedSet* τ)
| *update-element-type* (*Bag* -) τ (*Bag* τ)
| *update-element-type* (*Sequence* -) τ (*Sequence* τ)

inductive *to-unique-collection* **where**
to-unique-collection (*Collection* τ) (*Set* τ)
| *to-unique-collection* (*Set* τ) (*Set* τ)
| *to-unique-collection* (*OrderedSet* τ) (*OrderedSet* τ)
| *to-unique-collection* (*Bag* τ) (*Set* τ)
| *to-unique-collection* (*Sequence* τ) (*OrderedSet* τ)

inductive *to-nonunique-collection* **where**
to-nonunique-collection (*Collection* τ) (*Bag* τ)
| *to-nonunique-collection* (*Set* τ) (*Bag* τ)
| *to-nonunique-collection* (*OrderedSet* τ) (*Sequence* τ)
| *to-nonunique-collection* (*Bag* τ) (*Bag* τ)
| *to-nonunique-collection* (*Sequence* τ) (*Sequence* τ)

inductive *to-ordered-collection* **where**
to-ordered-collection (*Collection* τ) (*Sequence* τ)
| *to-ordered-collection* (*Set* τ) (*OrderedSet* τ)
| *to-ordered-collection* (*OrderedSet* τ) (*OrderedSet* τ)
| *to-ordered-collection* (*Bag* τ) (*Sequence* τ)
| *to-ordered-collection* (*Sequence* τ) (*Sequence* τ)

fun *to-single-type* **where**
to-single-type OclSuper = *OclSuper*
| *to-single-type* $\tau[1]$ = $\tau[1]$
| *to-single-type* $\tau[?]$ = $\tau[?]$
| *to-single-type* (*Collection* τ) = *to-single-type* τ
| *to-single-type* (*Set* τ) = *to-single-type* τ
| *to-single-type* (*OrderedSet* τ) = *to-single-type* τ
| *to-single-type* (*Bag* τ) = *to-single-type* τ
| *to-single-type* (*Sequence* τ) = *to-single-type* τ
| *to-single-type* (*Tuple* π) = *Tuple* π

fun *to-required-type* **where**
to-required-type $\tau[1]$ = $\tau[1]$
| *to-required-type* $\tau[?]$ = $\tau[1]$
| *to-required-type* τ = τ

fun *to-optional-type-nested* **where**
to-optional-type-nested OclSuper = *OclSuper*
| *to-optional-type-nested* $\tau[1]$ = $\tau[?]$

| *to-optional-type-nested* $\tau[?] = \tau[?]$
| *to-optional-type-nested* (*Collection* τ) = *Collection* (*to-optional-type-nested* τ)
| *to-optional-type-nested* (*Set* τ) = *Set* (*to-optional-type-nested* τ)
| *to-optional-type-nested* (*OrderedSet* τ) = *OrderedSet* (*to-optional-type-nested* τ)

| *to-optional-type-nested* (*Bag* τ) = *Bag* (*to-optional-type-nested* τ)
| *to-optional-type-nested* (*Sequence* τ) = *Sequence* (*to-optional-type-nested* τ)
| *to-optional-type-nested* (*Tuple* π) = *Tuple* (*fmmap to-optional-type-nested* π)

3.6 Determinism

lemma *element-type-det*:
element-type τ σ_1 $\Longrightarrow$
element-type τ σ_2 $\Longrightarrow$ $\sigma_1 = \sigma_2$
by (*induct rule*: *element-type.induct*; *simp add*: *element-type.simps*)

lemma *update-element-type-det*:
update-element-type τ σ ϱ_1 $\Longrightarrow$
update-element-type τ σ ϱ_2 $\Longrightarrow$ $\varrho_1 = \varrho_2$
by (*induct rule*: *update-element-type.induct*; *simp add*: *update-element-type.simps*)

lemma *to-unique-collection-det*:
to-unique-collection τ σ_1 $\Longrightarrow$
to-unique-collection τ σ_2 $\Longrightarrow$ $\sigma_1 = \sigma_2$
by (*induct rule*: *to-unique-collection.induct*; *simp add*: *to-unique-collection.simps*)

lemma *to-nonunique-collection-det*:
to-nonunique-collection τ σ_1 $\Longrightarrow$
to-nonunique-collection τ σ_2 $\Longrightarrow$ $\sigma_1 = \sigma_2$
by (*induct rule*: *to-nonunique-collection.induct*; *simp add*: *to-nonunique-collection.simps*)

lemma *to-ordered-collection-det*:
to-ordered-collection τ σ_1 $\Longrightarrow$
to-ordered-collection τ σ_2 $\Longrightarrow$ $\sigma_1 = \sigma_2$
by (*induct rule*: *to-ordered-collection.induct*; *simp add*: *to-ordered-collection.simps*)

3.7 Code Setup

code-pred *subtype* **.**

function *subtype-fun* :: *'a*::*order type* $\Rightarrow$ *'a type* $\Rightarrow$ *bool* **where**
subtype-fun OclSuper - = *False*
| *subtype-fun* (*Required* τ) σ = (*case* σ
of OclSuper $\Rightarrow$ *True*
| $\varrho[1]$ $\Rightarrow$ *basic-subtype-fun* τ ϱ
| $\varrho[?]$ $\Rightarrow$ *basic-subtype-fun* τ ϱ $\vee$ $\tau = \varrho$
| - $\Rightarrow$ *False*)
| *subtype-fun* (*Optional* τ) σ = (*case* σ

```
    of OclSuper ⇒ True
     | ϱ[?] ⇒ basic-subtype-fun τ ϱ
     | - ⇒ False)
|  subtype-fun (Collection τ) σ = (case σ
    of OclSuper ⇒ True
     | Collection ϱ ⇒ subtype-fun τ ϱ
     | - ⇒ False)
|  subtype-fun (Set τ) σ = (case σ
    of OclSuper ⇒ True
     | Collection ϱ ⇒ subtype-fun τ ϱ ∨ τ = ϱ
     | Set ϱ ⇒ subtype-fun τ ϱ
     | - ⇒ False)
|  subtype-fun (OrderedSet τ) σ = (case σ
    of OclSuper ⇒ True
     | Collection ϱ ⇒ subtype-fun τ ϱ ∨ τ = ϱ
     | OrderedSet ϱ ⇒ subtype-fun τ ϱ
     | - ⇒ False)
|  subtype-fun (Bag τ) σ = (case σ
    of OclSuper ⇒ True
     | Collection ϱ ⇒ subtype-fun τ ϱ ∨ τ = ϱ
     | Bag ϱ ⇒ subtype-fun τ ϱ
     | - ⇒ False)
|  subtype-fun (Sequence τ) σ = (case σ
    of OclSuper ⇒ True
     | Collection ϱ ⇒ subtype-fun τ ϱ ∨ τ = ϱ
     | Sequence ϱ ⇒ subtype-fun τ ϱ
     | - ⇒ False)
|  subtype-fun (Tuple π) σ = (case σ
    of OclSuper ⇒ True
     | Tuple ξ ⇒ strict-subtuple-fun (λτ σ. subtype-fun τ σ ∨ τ = σ) π ξ
     | - ⇒ False)
  by pat-completeness auto
termination
  by (relation  measure (λ(xs, ys). size ys) ;
      auto simp add: elem-le-ffold′ fmran′I)

lemma less-type-code [code]:
   (<) = subtype-fun
proof (intro ext iffI)
  fix τ σ ::  'a type
  show  τ < σ ⟹ subtype-fun τ σ
  proof (induct τ arbitrary: σ)
    case OclSuper thus ?case by (cases σ; auto)
  next
    case (Required τ) thus ?case
      by (cases σ; auto simp: less-basic-type-code less-eq-basic-type-code)
  next
    case (Optional τ) thus ?case
      by (cases σ; auto simp: less-basic-type-code less-eq-basic-type-code)
```

next
case (*Collection* τ) **thus** *?case* **by** (*cases* σ; *auto*)
next
case (*Set* τ) **thus** *?case* **by** (*cases* σ; *auto*)
next
case (*OrderedSet* τ) **thus** *?case* **by** (*cases* σ; *auto*)
next
case (*Bag* τ) **thus** *?case* **by** (*cases* σ; *auto*)
next
case (*Sequence* τ) **thus** *?case* **by** (*cases* σ; *auto*)
next
case (*Tuple* π)
have
$\bigwedge\xi$. *subtuple* $(\leq)$ π ξ $\longrightarrow$
subtuple $(\lambda\tau$ σ. *subtype-fun* τ $\sigma \vee \tau = \sigma)$ π ξ
by (*rule subtuple-mono*; *auto simp add*: *Tuple.hyps*)
with *Tuple.prems* **show** *?case* **by** (*cases* σ; *auto*)
qed
show *subtype-fun* τ $\sigma \Longrightarrow \tau < \sigma$
proof (*induct* σ *arbitrary*: τ)
case *OclSuper* **thus** *?case* **by** (*cases* σ; *auto*)
next
case (*Required* σ) **show** *?case*
by (*insert Required*) (*erule subtype-fun.elims*;
auto simp: *less-basic-type-code less-eq-basic-type-code*)
next
case (*Optional* σ) **show** *?case*
by (*insert Optional*) (*erule subtype-fun.elims*;
auto simp: *less-basic-type-code less-eq-basic-type-code*)
next
case (*Collection* σ) **show** *?case*
apply (*insert Collection*)
apply (*erule subtype-fun.elims*; *auto*)
using *order.strict-implies-order* **by** *auto*
next
case (*Set* σ) **show** *?case*
by (*insert Set*) (*erule subtype-fun.elims*; *auto*)
next
case (*OrderedSet* σ) **show** *?case*
by (*insert OrderedSet*) (*erule subtype-fun.elims*; *auto*)
next
case (*Bag* σ) **show** *?case*
by (*insert Bag*) (*erule subtype-fun.elims*; *auto*)
next
case (*Sequence* σ) **show** *?case*
by (*insert Sequence*) (*erule subtype-fun.elims*; *auto*)
next
case (*Tuple* ξ)
have *subtuple-imp-simp*:

```
      ⋀π. subtuple (λτ σ. subtype-fun τ σ ∨ τ = σ) π ξ ⟶
      subtuple (≤) π ξ
     by (rule subtuple-mono; auto simp add: Tuple.hyps less-imp-le)
   show ?case
     apply (insert Tuple)
     by (erule subtype-fun.elims; auto simp add: subtuple-imp-simp)
 qed
qed

lemma less-eq-type-code [code]:
  (≤) = (λx y. subtype-fun x y ∨ x = y)
 unfolding dual-order.order-iff-strict less-type-code
 by auto

code-pred element-type .
code-pred update-element-type .
code-pred to-unique-collection .
code-pred to-nonunique-collection .
code-pred to-ordered-collection .

end
```

Chapter 4

Abstract Syntax

theory *OCL-Syntax*
imports *Complex-Main Object-Model OCL-Types*
begin

4.1 Preliminaries

type-synonym *vname* = *String.literal*
type-synonym *'a env* = *vname* $\rightharpoonup_f$ *'a*

In OCL $1 + \infty = \bot$. So we do not use *enat* and define the new data type.

typedef *unat* = *UNIV* :: *nat option set* **..**

definition *unat x* ≡ *Abs-unat* (*Some x*)

instantiation *unat* :: *infinity*
begin
definition $\infty \equiv$ *Abs-unat None*
instance ..
end

free-constructors *cases-unat* **for**
unat
| ∞ :: *unat*
unfolding *unat-def infinity-unat-def*
apply (*metis Rep-unat-inverse option.collapse*)
apply (*metis Abs-unat-inverse UNIV-I option.sel*)
by (*simp add*: *Abs-unat-inject*)

4.2 Standard Library Operations

datatype *metaop* = *AllInstancesOp*

datatype *typeop* = *OclAsTypeOp* | *OclIsTypeOfOp* | *OclIsKindOfOp*
| *SelectByKindOp* | *SelectByTypeOp*

datatype *super-binop* = *EqualOp* | *NotEqualOp*

datatype *any-unop* = *OclAsSetOp* | *OclIsNewOp*
| *OclIsUndefinedOp* | *OclIsInvalidOp* | *OclLocaleOp* | *ToStringOp*

datatype *boolean-unop* = *NotOp*
datatype *boolean-binop* = *AndOp* | *OrOp* | *XorOp* | *ImpliesOp*

datatype *numeric-unop* = *UMinusOp* | *AbsOp* | *FloorOp* | *RoundOp* | *ToIntegerOp*
datatype *numeric-binop* = *PlusOp* | *MinusOp* | *MultOp* | *DivideOp*
| *DivOp* | *ModOp* | *MaxOp* | *MinOp*
| *LessOp* | *LessEqOp* | *GreaterOp* | *GreaterEqOp*

datatype *string-unop* = *SizeOp* | *ToUpperCaseOp* | *ToLowerCaseOp* | *Character-sOp*
| *ToBooleanOp* | *ToIntegerOp* | *ToRealOp*
datatype *string-binop* = *ConcatOp* | *IndexOfOp* | *EqualsIgnoreCaseOp* | *AtOp*
| *LessOp* | *LessEqOp* | *GreaterOp* | *GreaterEqOp*
datatype *string-ternop* = *SubstringOp*

datatype *collection-unop* = *CollectionSizeOp* | *IsEmptyOp* | *NotEmptyOp*
| *CollectionMaxOp* | *CollectionMinOp* | *SumOp*
| *AsSetOp* | *AsOrderedSetOp* | *AsSequenceOp* | *AsBagOp* | *FlattenOp*
| *FirstOp* | *LastOp* | *ReverseOp*
datatype *collection-binop* = *IncludesOp* | *ExcludesOp*
| *CountOp*| *IncludesAllOp* | *ExcludesAllOp* | *ProductOp*
| *UnionOp* | *IntersectionOp* | *SetMinusOp* | *SymmetricDifferenceOp*
| *IncludingOp* | *ExcludingOp*
| *AppendOp* | *PrependOp* | *CollectionAtOp* | *CollectionIndexOfOp*
datatype *collection-ternop* = *InsertAtOp* | *SubOrderedSetOp* | *SubSequenceOp*

type-synonym *unop* = *any-unop* + *boolean-unop* + *numeric-unop* + *string-unop* + *collection-unop*

declare [[*coercion* *Inl* :: *any-unop* ⇒ *unop*]]
declare [[*coercion* *Inr* ∘ *Inl* :: *boolean-unop* ⇒ *unop*]]
declare [[*coercion* *Inr* ∘ *Inr* ∘ *Inl* :: *numeric-unop* ⇒ *unop*]]
declare [[*coercion* *Inr* ∘ *Inr* ∘ *Inr* ∘ *Inl* :: *string-unop* ⇒ *unop*]]
declare [[*coercion* *Inr* ∘ *Inr* ∘ *Inr* ∘ *Inr* :: *collection-unop* ⇒ *unop*]]

type-synonym *binop* = *super-binop* + *boolean-binop* + *numeric-binop* + *string-binop* + *collection-binop*

declare [[*coercion* *Inl* :: *super-binop* ⇒ *binop*]]
declare [[*coercion* *Inr* ∘ *Inl* :: *boolean-binop* ⇒ *binop*]]
declare [[*coercion* *Inr* ∘ *Inr* ∘ *Inl* :: *numeric-binop* ⇒ *binop*]]

declare [[*coercion Inr* ∘ *Inr* ∘ *Inr* ∘ *Inl* :: *string-binop* ⇒ *binop*]]
declare [[*coercion Inr* ∘ *Inr* ∘ *Inr* ∘ *Inr* :: *collection-binop* ⇒ *binop*]]

type-synonym *ternop* = *string-ternop* + *collection-ternop*

declare [[*coercion Inl* :: *string-ternop* ⇒ *ternop*]]
declare [[*coercion Inr* :: *collection-ternop* ⇒ *ternop*]]

type-synonym *op* = *unop* + *binop* + *ternop* + *oper*

declare [[*coercion Inl* ∘ *Inl* :: *any-unop* ⇒ *op*]]
declare [[*coercion Inl* ∘ *Inr* ∘ *Inl* :: *boolean-unop* ⇒ *op*]]
declare [[*coercion Inl* ∘ *Inr* ∘ *Inr* ∘ *Inl* :: *numeric-unop* ⇒ *op*]]
declare [[*coercion Inl* ∘ *Inr* ∘ *Inr* ∘ *Inr* ∘ *Inl* :: *string-unop* ⇒ *op*]]
declare [[*coercion Inl* ∘ *Inr* ∘ *Inr* ∘ *Inr* ∘ *Inr* :: *collection-unop* ⇒ *op*]]

declare [[*coercion Inr* ∘ *Inl* ∘ *Inl* :: *super-binop* ⇒ *op*]]
declare [[*coercion Inr* ∘ *Inl* ∘ *Inr* ∘ *Inl* :: *boolean-binop* ⇒ *op*]]
declare [[*coercion Inr* ∘ *Inl* ∘ *Inr* ∘ *Inr* ∘ *Inl* :: *numeric-binop* ⇒ *op*]]
declare [[*coercion Inr* ∘ *Inl* ∘ *Inr* ∘ *Inr* ∘ *Inr* ∘ *Inl* :: *string-binop* ⇒ *op*]]
declare [[*coercion Inr* ∘ *Inl* ∘ *Inr* ∘ *Inr* ∘ *Inr* ∘ *Inr* :: *collection-binop* ⇒ *op*]]

declare [[*coercion Inr* ∘ *Inr* ∘ *Inl* ∘ *Inl* :: *string-ternop* ⇒ *op*]]
declare [[*coercion Inr* ∘ *Inr* ∘ *Inl* ∘ *Inr* :: *collection-ternop* ⇒ *op*]]

declare [[*coercion Inr* ∘ *Inr* ∘ *Inr* :: *oper* ⇒ *op*]]

datatype *iterator* = *AnyIter* | *ClosureIter* | *CollectIter* | *CollectNestedIter*
| *ExistsIter* | *ForAllIter* | *OneIter* | *IsUniqueIter*
| *SelectIter* | *RejectIter* | *SortedByIter*

4.3 Expressions

datatype *collection-literal-kind* =
CollectionKind | *SetKind* | *OrderedSetKind* | *BagKind* | *SequenceKind*

A call kind could be defined as two boolean values (*is-arrow-call*, *is-safe-call*). Also we could derive *is-arrow-call* value automatically based on an operation kind. However, it is much easier and more natural to use the following enumeration.

datatype *call-kind* = *DotCall* | *ArrowCall* | *SafeDotCall* | *SafeArrowCall*

We do not define a *Classifier* type (a type of all types), because it will add unnecessary complications to the theory. So we have to define type operations as a pure syntactic constructs. We do not define *Type* expressions either.

We do not define *InvalidLiteral*, because it allows us to exclude *OclInvalid* type from typing rules. It simplifies the types system.

Please take a note that for *AssociationEnd* and *AssociationClass* call expressions one can specify an optional role of a source class (*from-role*). It differs from the OCL specification, which allows one to specify a role of a destination class. However, the latter one does not allow one to determine uniquely a set of linked objects, for example, in a ternary self relation.

datatype *'a expr* =
  Literal 'a literal-expr
| *Let* (*var* : *vname*) (*var-type* : *'a type option*) (*init-expr* : *'a expr*)
    (*body-expr* : *'a expr*)
| *Var* (*var* : *vname*)
| *If* (*if-expr* : *'a expr*) (*then-expr* : *'a expr*) (*else-expr* : *'a expr*)
| *MetaOperationCall* (*type* : *'a type*) *metaop*
| *StaticOperationCall* (*type* : *'a type*) *oper* (*args* : *'a expr list*)
| *Call* (*source* : *'a expr*) (*kind* : *call-kind*) *'a call-expr*
and *'a literal-expr* =
  NullLiteral
| *BooleanLiteral* (*boolean-symbol* : *bool*)
| *RealLiteral* (*real-symbol* : *real*)
| *IntegerLiteral* (*integer-symbol* : *int*)
| *UnlimitedNaturalLiteral* (*unlimited-natural-symbol* : *unat*)
| *StringLiteral* (*string-symbol* : *string*)
| *EnumLiteral* (*enum-type* : *'a enum*) (*enum-literal* : *elit*)
| *CollectionLiteral* (*kind* : *collection-literal-kind*)
    (*parts* : *'a collection-literal-part-expr list*)
| *TupleLiteral* (*tuple-elements* : (*telem* × *'a type option* × *'a expr*) *list*)
and *'a collection-literal-part-expr* =
  CollectionItem (*item* : *'a expr*)
| *CollectionRange* (*first* : *'a expr*) (*last* : *'a expr*)
and *'a call-expr* =
  TypeOperation typeop (*type* : *'a type*)
| *Attribute attr*
| *AssociationEnd* (*from-role* : *role option*) *role*
| *AssociationClass* (*from-role* : *role option*) *'a*
| *AssociationClassEnd role*
| *Operation op* (*args* : *'a expr list*)
| *TupleElement telem*
| *Iterate* (*iterators* : *vname list*) (*iterators-type* : *'a type option*)
    (*var* : *vname*) (*var-type* : *'a type option*) (*init-expr* : *'a expr*)
    (*body-expr* : *'a expr*)
| *Iterator iterator* (*iterators* : *vname list*) (*iterators-type* : *'a type option*)
    (*body-expr* : *'a expr*)

definition *tuple-element-name* ≡ *fst*
definition *tuple-element-type* ≡ *fst* ∘ *snd*
definition *tuple-element-expr* ≡ *snd* ∘ *snd*

declare [[*coercion Literal* :: *'a literal-expr* ⇒ *'a expr*]]

abbreviation *TypeOperationCall src k op ty* ≡
Call src k (*TypeOperation op ty*)
abbreviation *AttributeCall src k attr* ≡
Call src k (*Attribute attr*)
abbreviation *AssociationEndCall src k from role* ≡
Call src k (*AssociationEnd from role*)
abbreviation *AssociationClassCall src k from cls* ≡
Call src k (*AssociationClass from cls*)
abbreviation *AssociationClassEndCall src k role* ≡
Call src k (*AssociationClassEnd role*)
abbreviation *OperationCall src k op as* ≡
Call src k (*Operation op as*)
abbreviation *TupleElementCall src k elem* ≡
Call src k (*TupleElement elem*)
abbreviation *IterateCall src k its its-ty v ty init body* ≡
Call src k (*Iterate its its-ty v ty init body*)
abbreviation *AnyIteratorCall src k its its-ty body* ≡
Call src k (*Iterator AnyIter its its-ty body*)
abbreviation *ClosureIteratorCall src k its its-ty body* ≡
Call src k (*Iterator ClosureIter its its-ty body*)
abbreviation *CollectIteratorCall src k its its-ty body* ≡
Call src k (*Iterator CollectIter its its-ty body*)
abbreviation *CollectNestedIteratorCall src k its its-ty body* ≡
Call src k (*Iterator CollectNestedIter its its-ty body*)
abbreviation *ExistsIteratorCall src k its its-ty body* ≡
Call src k (*Iterator ExistsIter its its-ty body*)
abbreviation *ForAllIteratorCall src k its its-ty body* ≡
Call src k (*Iterator ForAllIter its its-ty body*)
abbreviation *OneIteratorCall src k its its-ty body* ≡
Call src k (*Iterator OneIter its its-ty body*)
abbreviation *IsUniqueIteratorCall src k its its-ty body* ≡
Call src k (*Iterator IsUniqueIter its its-ty body*)
abbreviation *SelectIteratorCall src k its its-ty body* ≡
Call src k (*Iterator SelectIter its its-ty body*)
abbreviation *RejectIteratorCall src k its its-ty body* ≡
Call src k (*Iterator RejectIter its its-ty body*)
abbreviation *SortedByIteratorCall src k its its-ty body* ≡
Call src k (*Iterator SortedByIter its its-ty body*)

end

Chapter 5

Object Model

theory *OCL-Object-Model*
imports *OCL-Syntax*
begin

I see no reason why objects should refer nulls using multi-valued associations. Therefore, multi-valued associations have collection types with non-nullable element types.

definition
assoc-end-type end ≡
let $\mathcal{C}$ = *assoc-end-class end in*
if assoc-end-max end ≤ (*1* :: *nat*) *then*
if assoc-end-min end = (*0* :: *nat*)
then $\langle\mathcal{C}\rangle_{\mathcal{T}}[?]$
else $\langle\mathcal{C}\rangle_{\mathcal{T}}[1]$
else
if assoc-end-unique end then
if assoc-end-ordered end
then OrderedSet $\langle\mathcal{C}\rangle_{\mathcal{T}}[1]$
else Set $\langle\mathcal{C}\rangle_{\mathcal{T}}[1]$
else
if assoc-end-ordered end
then Sequence $\langle\mathcal{C}\rangle_{\mathcal{T}}[1]$
else Bag $\langle\mathcal{C}\rangle_{\mathcal{T}}[1]$

definition *class-assoc-type* $\mathcal{A}$ ≡ *Set* $\langle\mathcal{A}\rangle_{\mathcal{T}}[1]$

definition *class-assoc-end-type end* ≡ $\langle$*assoc-end-class end*$\rangle_{\mathcal{T}}[1]$

definition *oper-type op* ≡
let params = *oper-out-params op in*
if length params = *0*
then oper-result op
else Tuple (*fmap-of-list* (*map* (λ*p*. (*param-name p*, *param-type p*))
(*params* @ [(*STR* ''*result*'', *oper-result op*, *Out*)])))

class *ocl-object-model* =
fixes *classes* :: *'a* :: *semilattice-sup fset*
and *attributes* :: *'a* $\rightharpoonup_f$ *attr* $\rightharpoonup_f$ *'a type*
and *associations* :: *assoc* $\rightharpoonup_f$ *role* $\rightharpoonup_f$ *'a assoc-end*
and *association-classes* :: *'a* $\rightharpoonup_f$ *assoc*
and *operations* :: (*'a type*, *'a expr*) *oper-spec list*
and *literals* :: *'a enum* $\rightharpoonup_f$ *elit fset*
assumes *assoc-end-min-less-eq-max*:
assoc $|\in|$ *fmdom associations* $\Longrightarrow$
fmlookup associations assoc = *Some ends* $\Longrightarrow$
role $|\in|$ *fmdom ends* $\Longrightarrow$
fmlookup ends role = *Some end* $\Longrightarrow$
assoc-end-min end $\leq$ *assoc-end-max end*
assumes *association-ends-unique*:
association-ends' *classes associations* $\mathcal{C}$ *from role* end_1 $\Longrightarrow$
association-ends' *classes associations* $\mathcal{C}$ *from role* end_2 $\Longrightarrow$ $end_1 = end_2$
begin

interpretation *base*: *object-model*
by *standard* (*simp-all add*: *local.assoc-end-min-less-eq-max local.association-ends-unique*)

abbreviation *owned-attribute* $\equiv$ *base.owned-attribute*
abbreviation *attribute* $\equiv$ *base.attribute*
abbreviation *association-ends* $\equiv$ *base.association-ends*
abbreviation *owned-association-end* $\equiv$ *base.owned-association-end*
abbreviation *association-end* $\equiv$ *base.association-end*
abbreviation *referred-by-association-class* $\equiv$ *base.referred-by-association-class*
abbreviation *association-class-end* $\equiv$ *base.association-class-end*
abbreviation *operation* $\equiv$ *base.operation*
abbreviation *operation-defined* $\equiv$ *base.operation-defined*
abbreviation *static-operation* $\equiv$ *base.static-operation*
abbreviation *static-operation-defined* $\equiv$ *base.static-operation-defined*
abbreviation *has-literal* $\equiv$ *base.has-literal*

lemmas *attribute-det* = *base.attribute-det*
lemmas *attribute-self-or-inherited* = *base.attribute-self-or-inherited*
lemmas *attribute-closest* = *base.attribute-closest*
lemmas *association-end-det* = *base.association-end-det*
lemmas *association-end-self-or-inherited* = *base.association-end-self-or-inherited*
lemmas *association-end-closest* = *base.association-end-closest*
lemmas *association-class-end-det* = *base.association-class-end-det*
lemmas *operation-det* = *base.operation-det*
lemmas *static-operation-det* = *base.static-operation-det*

end

end

Chapter 6

Typing

theory *OCL-Typing*
imports *OCL-Object-Model HOL−Library.Transitive-Closure-Table*
begin

The following rules are more restrictive than rules given in the OCL specification. This allows one to identify more errors in expressions. However, these restrictions may be revised if necessary. Perhaps some of them could be separated and should cause warnings instead of errors.

6.1 Operations Typing

6.1.1 Metaclass Operations

All basic types in the theory are either nullable or non-nullable. For example, instead of *Boolean* type we have two types: *Boolean*[*1*] and *Boolean*[*?*]. The *allInstances*() operation is extended accordingly:

```
Boolean[1].allInstances() = Set{true, false}
Boolean[?].allInstances() = Set{true, false, null}
```

inductive *mataop-type* **where**
mataop-type τ *AllInstancesOp* (*Set* τ)

6.1.2 Type Operations

At first we decided to allow casting only to subtypes. However sometimes it is necessary to cast expressions to supertypes, for example, to access overridden attributes of a supertype. So we allow casting to subtypes and supertypes. Casting to other types is meaningless.

According to the Section 7.4.7 of the OCL specification *oclAsType*() can be applied to collections as well as to single values. I guess we can allow *oclIsTypeOf*() and *oclIsKindOf*() for collections too.

Please take a note that the following expressions are prohibited, because they always return true or false:

```
1.oclIsKindOf(OclAny[?])
1.oclIsKindOf(String[1])
```

Please take a note that:

```
Set{1,2,null,'abc'}->selectByKind(Integer[1]) = Set{1,2}
Set{1,2,null,'abc'}->selectByKind(Integer[?]) = Set{1,2,null}
```

The following expressions are prohibited, because they always returns either the same or empty collections:

```
Set{1,2,null,'abc'}->selectByKind(OclAny[?])
Set{1,2,null,'abc'}->selectByKind(Collection(Boolean[1]))
```

inductive *typeop-type* **where**
$\sigma < \tau \vee \tau < \sigma \Longrightarrow$
typeop-type DotCall OclAsTypeOp $\tau\ \sigma\ \sigma$

| $\sigma < \tau \Longrightarrow$
typeop-type DotCall OclIsTypeOfOp $\tau\ \sigma$ *Boolean*[*1*]
| $\sigma < \tau \Longrightarrow$
typeop-type DotCall OclIsKindOfOp $\tau\ \sigma$ *Boolean*[*1*]

| *element-type* $\tau\ \varrho \Longrightarrow \sigma < \varrho \Longrightarrow$
update-element-type $\tau\ \sigma\ \upsilon \Longrightarrow$
typeop-type ArrowCall SelectByKindOp $\tau\ \sigma\ \upsilon$

| *element-type* $\tau\ \varrho \Longrightarrow \sigma < \varrho \Longrightarrow$
update-element-type $\tau\ \sigma\ \upsilon \Longrightarrow$
typeop-type ArrowCall SelectByTypeOp $\tau\ \sigma\ \upsilon$

6.1.3 OclSuper Operations

It makes sense to compare values only with compatible types.

inductive *super-binop-type*
:: *super-binop* $\Rightarrow$ ('*a* :: *order*) *type* $\Rightarrow$ '*a type* $\Rightarrow$ '*a type* $\Rightarrow$ *bool* **where**
$\tau \leq \sigma \vee \sigma < \tau \Longrightarrow$
super-binop-type EqualOp $\tau\ \sigma$ *Boolean*[*1*]
| $\tau \leq \sigma \vee \sigma < \tau \Longrightarrow$
super-binop-type NotEqualOp $\tau\ \sigma$ *Boolean*[*1*]

6.1.4 OclAny Operations

The OCL specification defines *toString*() operation only for boolean and numeric types. However, I guess it is a good idea to define it once for all basic types. Maybe it should be defined for collections as well.

inductive *any-unop-type* **where**

$\tau \leq OclAny[?] \Longrightarrow$
any-unop-type OclAsSetOp τ (*Set* (*to-required-type* τ))
| $\tau \leq OclAny[?] \Longrightarrow$
any-unop-type OclIsNewOp τ *Boolean*[*1*]
| $\tau \leq OclAny[?] \Longrightarrow$
any-unop-type OclIsUndefinedOp τ *Boolean*[*1*]
| $\tau \leq OclAny[?] \Longrightarrow$
any-unop-type OclIsInvalidOp τ *Boolean*[*1*]
| $\tau \leq OclAny[?] \Longrightarrow$
any-unop-type OclLocaleOp τ *String*[*1*]
| $\tau \leq OclAny[?] \Longrightarrow$
any-unop-type ToStringOp τ *String*[*1*]

6.1.5 Boolean Operations

Please take a note that:

```
true or false : Boolean[1]
true and null : Boolean[?]
null and null : OclVoid[?]
```

inductive *boolean-unop-type* **where**
$\tau \leq Boolean[?] \Longrightarrow$
boolean-unop-type NotOp τ τ

inductive *boolean-binop-type* **where**
$\tau \sqcup \sigma = \varrho \Longrightarrow \varrho \leq Boolean[?] \Longrightarrow$
boolean-binop-type AndOp τ σ ϱ
| $\tau \sqcup \sigma = \varrho \Longrightarrow \varrho \leq Boolean[?] \Longrightarrow$
boolean-binop-type OrOp τ σ ϱ
| $\tau \sqcup \sigma = \varrho \Longrightarrow \varrho \leq Boolean[?] \Longrightarrow$
boolean-binop-type XorOp τ σ ϱ
| $\tau \sqcup \sigma = \varrho \Longrightarrow \varrho \leq Boolean[?] \Longrightarrow$
boolean-binop-type ImpliesOp τ σ ϱ

6.1.6 Numeric Operations

The expression *1 + null* is not well-typed. Nullable numeric values should be converted to non-nullable ones. This is a significant difference from the OCL specification.

Please take a note that many operations automatically casts unlimited naturals to integers.

The difference between *oclAsType*(*Integer*) and *toInteger*() for unlimited naturals is unclear.

inductive *numeric-unop-type* **where**
$\tau = Real[1] \Longrightarrow$
numeric-unop-type UMinusOp τ *Real*[*1*]
| $\tau = UnlimitedNatural[1] - Integer[1] \Longrightarrow$

$\mathit{numeric\text{-}unop\text{-}type}\ \mathit{UMinusOp}\ \tau\ \mathit{Integer}[1]$

$\mid\ \tau = \mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}unop\text{-}type}\ \mathit{AbsOp}\ \tau\ \mathit{Real}[1]$
$\mid\ \tau = \mathit{UnlimitedNatural}[1]{-}\mathit{Integer}[1] \Longrightarrow$
$\mathit{numeric\text{-}unop\text{-}type}\ \mathit{AbsOp}\ \tau\ \mathit{Integer}[1]$

$\mid\ \tau = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}unop\text{-}type}\ \mathit{FloorOp}\ \tau\ \mathit{Integer}[1]$
$\mid\ \tau = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}unop\text{-}type}\ \mathit{RoundOp}\ \tau\ \mathit{Integer}[1]$

$\mid\ \tau = \mathit{UnlimitedNatural}[1] \Longrightarrow$
$\mathit{numeric\text{-}unop\text{-}type}\ \mathit{numeric\text{-}unop.ToIntegerOp}\ \tau\ \mathit{Integer}[1]$

inductive *numeric-binop-type* **where**
$\tau \sqcup \sigma = \varrho \Longrightarrow \varrho = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{PlusOp}\ \tau\ \sigma\ \varrho$

$\mid\ \tau \sqcup \sigma = \mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{MinusOp}\ \tau\ \sigma\ \mathit{Real}[1]$
$\mid\ \tau \sqcup \sigma = \mathit{UnlimitedNatural}[1]{-}\mathit{Integer}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{MinusOp}\ \tau\ \sigma\ \mathit{Integer}[1]$

$\mid\ \tau \sqcup \sigma = \varrho \Longrightarrow \varrho = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{MultOp}\ \tau\ \sigma\ \varrho$

$\mid\ \tau = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow \sigma = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{DivideOp}\ \tau\ \sigma\ \mathit{Real}[1]$

$\mid\ \tau \sqcup \sigma = \varrho \Longrightarrow \varrho = \mathit{UnlimitedNatural}[1]{-}\mathit{Integer}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{DivOp}\ \tau\ \sigma\ \varrho$
$\mid\ \tau \sqcup \sigma = \varrho \Longrightarrow \varrho = \mathit{UnlimitedNatural}[1]{-}\mathit{Integer}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{ModOp}\ \tau\ \sigma\ \varrho$

$\mid\ \tau \sqcup \sigma = \varrho \Longrightarrow \varrho = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{MaxOp}\ \tau\ \sigma\ \varrho$
$\mid\ \tau \sqcup \sigma = \varrho \Longrightarrow \varrho = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{MinOp}\ \tau\ \sigma\ \varrho$

$\mid\ \tau = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow \sigma = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{numeric\text{-}binop.LessOp}\ \tau\ \sigma\ \mathit{Boolean}[1]$
$\mid\ \tau = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow \sigma = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{numeric\text{-}binop.LessEqOp}\ \tau\ \sigma\ \mathit{Boolean}[1]$
$\mid\ \tau = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow \sigma = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{numeric\text{-}binop.GreaterOp}\ \tau\ \sigma\ \mathit{Boolean}[1]$
$\mid\ \tau = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow \sigma = \mathit{UnlimitedNatural}[1]{-}\mathit{Real}[1] \Longrightarrow$
$\mathit{numeric\text{-}binop\text{-}type}\ \mathit{numeric\text{-}binop.GreaterEqOp}\ \tau\ \sigma\ \mathit{Boolean}[1]$

6.1.7 String Operations

inductive *string-unop-type* **where**
string-unop-type SizeOp String[1] Integer[1]
| *string-unop-type CharactersOp String[1] (Sequence String[1])*
| *string-unop-type ToUpperCaseOp String[1] String[1]*
| *string-unop-type ToLowerCaseOp String[1] String[1]*
| *string-unop-type ToBooleanOp String[1] Boolean[1]*
| *string-unop-type ToIntegerOp String[1] Integer[1]*
| *string-unop-type ToRealOp String[1] Real[1]*

inductive *string-binop-type* **where**
string-binop-type ConcatOp String[1] String[1] String[1]
| *string-binop-type EqualsIgnoreCaseOp String[1] String[1] Boolean[1]*
| *string-binop-type LessOp String[1] String[1] Boolean[1]*
| *string-binop-type LessEqOp String[1] String[1] Boolean[1]*
| *string-binop-type GreaterOp String[1] String[1] Boolean[1]*
| *string-binop-type GreaterEqOp String[1] String[1] Boolean[1]*
| *string-binop-type IndexOfOp String[1] String[1] Integer[1]*
| τ = *UnlimitedNatural[1]*−*Integer[1]* $\Longrightarrow$
string-binop-type AtOp String[1] τ *String[1]*

inductive *string-ternop-type* **where**
σ = *UnlimitedNatural[1]*−*Integer[1]* $\Longrightarrow$
ϱ = *UnlimitedNatural[1]*−*Integer[1]* $\Longrightarrow$
string-ternop-type SubstringOp String[1] σ ϱ *String[1]*

6.1.8 Collection Operations

Please take a note, that *flatten*() preserves a collection kind.

inductive *collection-unop-type* **where**
element-type τ - $\Longrightarrow$
collection-unop-type CollectionSizeOp τ *Integer[1]*
| *element-type* τ - $\Longrightarrow$
collection-unop-type IsEmptyOp τ *Boolean[1]*
| *element-type* τ - $\Longrightarrow$
collection-unop-type NotEmptyOp τ *Boolean[1]*

| *element-type* τ σ $\Longrightarrow$ σ = *UnlimitedNatural[1]*−*Real[1]* $\Longrightarrow$
collection-unop-type CollectionMaxOp τ σ
| *element-type* τ σ $\Longrightarrow$ *operation* σ *STR ''max''* [σ] *oper* $\Longrightarrow$
collection-unop-type CollectionMaxOp τ σ

| *element-type* τ σ $\Longrightarrow$ σ = *UnlimitedNatural[1]*−*Real[1]* $\Longrightarrow$
collection-unop-type CollectionMinOp τ σ
| *element-type* τ σ $\Longrightarrow$ *operation* σ *STR ''min''* [σ] *oper* $\Longrightarrow$
collection-unop-type CollectionMinOp τ σ

| *element-type* τ σ $\Longrightarrow$ σ = *UnlimitedNatural[1]*−*Real[1]* $\Longrightarrow$

collection-unop-type SumOp τ σ
| *element-type* τ σ $\Longrightarrow$ *operation* σ *STR* ″+″ [σ] *oper* $\Longrightarrow$
collection-unop-type SumOp τ σ

| *element-type* τ σ $\Longrightarrow$
collection-unop-type AsSetOp τ (*Set* σ)
| *element-type* τ σ $\Longrightarrow$
collection-unop-type AsOrderedSetOp τ (*OrderedSet* σ)
| *element-type* τ σ $\Longrightarrow$
collection-unop-type AsBagOp τ (*Bag* σ)
| *element-type* τ σ $\Longrightarrow$
collection-unop-type AsSequenceOp τ (*Sequence* σ)

| *update-element-type* τ (*to-single-type* τ) σ $\Longrightarrow$
collection-unop-type FlattenOp τ σ

| *collection-unop-type FirstOp* (*OrderedSet* τ) τ
| *collection-unop-type FirstOp* (*Sequence* τ) τ
| *collection-unop-type LastOp* (*OrderedSet* τ) τ
| *collection-unop-type LastOp* (*Sequence* τ) τ
| *collection-unop-type ReverseOp* (*OrderedSet* τ) (*OrderedSet* τ)
| *collection-unop-type ReverseOp* (*Sequence* τ) (*Sequence* τ)

Please take a note that if both arguments are collections, then an element type of the resulting collection is a super type of element types of orginal collections. However for single-valued operations (*append*(), *insertAt*(), ...) this behavior looks undesirable. So we restrict such arguments to have a subtype of the collection element type.

Please take a note that we allow the following expressions:

```
let nullable_value : Integer[?] = null in
Sequence{1..3}->inculdes(nullable_value) and
Sequence{1..3}->inculdes(null) and
Sequence{1..3}->inculdesAll(Set{1,null})
```

The OCL specification defines *including*() and *excluding*() operations for the *Sequence* type but does not define them for the *OrderedSet* type. We define them for all collection types.

It is a good idea to prohibit including of values that do not conform to a collection element type. However we do not restrict it.

At first we defined the following typing rules for the *excluding*() operation:

| element-type τ ϱ $\Longrightarrow$ $\sigma \leq \varrho$ $\Longrightarrow$ $\sigma \neq$ OclVoid[?] $\Longrightarrow$
collection-binop-type ExcludingOp τ σ τ
| element-type τ ϱ $\Longrightarrow$ $\sigma \leq \varrho$ $\Longrightarrow$ $\sigma =$ OclVoid[?] $\Longrightarrow$
update-element-type τ (to-required-type ϱ) υ $\Longrightarrow$
collection-binop-type ExcludingOp τ σ υ

This operation could play a special role in a definition of safe navigation operations:

```
Sequence{1,2,null}->exculding(null) : Integer[1]
```

However it is more natural to use a *selectByKind*($T[1]$) operation instead.

inductive *collection-binop-type* **where**
element-type $\tau\ \varrho \Longrightarrow \sigma \leq$ *to-optional-type-nested* $\varrho \Longrightarrow$
collection-binop-type IncludesOp $\tau\ \sigma$ *Boolean*[*1*]
| *element-type* $\tau\ \varrho \Longrightarrow \sigma \leq$ *to-optional-type-nested* $\varrho \Longrightarrow$
collection-binop-type ExcludesOp $\tau\ \sigma$ *Boolean*[*1*]
| *element-type* $\tau\ \varrho \Longrightarrow \sigma \leq$ *to-optional-type-nested* $\varrho \Longrightarrow$
collection-binop-type CountOp $\tau\ \sigma$ *Integer*[*1*]

| *element-type* $\tau\ \varrho \Longrightarrow$ *element-type* $\sigma\ \upsilon \Longrightarrow \upsilon \leq$ *to-optional-type-nested* $\varrho \Longrightarrow$
collection-binop-type IncludesAllOp $\tau\ \sigma$ *Boolean*[*1*]
| *element-type* $\tau\ \varrho \Longrightarrow$ *element-type* $\sigma\ \upsilon \Longrightarrow \upsilon \leq$ *to-optional-type-nested* $\varrho \Longrightarrow$
collection-binop-type ExcludesAllOp $\tau\ \sigma$ *Boolean*[*1*]

| *element-type* $\tau\ \varrho \Longrightarrow$ *element-type* $\sigma\ \upsilon \Longrightarrow$
collection-binop-type ProductOp $\tau\ \sigma$
(*Set* (*Tuple* (*fmap-of-list* [(*STR* ''*first*'', ϱ), (*STR* ''*second*'', υ)])))

| *collection-binop-type UnionOp* (*Set* τ) (*Set* σ) (*Set* ($\tau \sqcup \sigma$))
| *collection-binop-type UnionOp* (*Set* τ) (*Bag* σ) (*Bag* ($\tau \sqcup \sigma$))
| *collection-binop-type UnionOp* (*Bag* τ) (*Set* σ) (*Bag* ($\tau \sqcup \sigma$))
| *collection-binop-type UnionOp* (*Bag* τ) (*Bag* σ) (*Bag* ($\tau \sqcup \sigma$))

| *collection-binop-type IntersectionOp* (*Set* τ) (*Set* σ) (*Set* ($\tau \sqcup \sigma$))
| *collection-binop-type IntersectionOp* (*Set* τ) (*Bag* σ) (*Set* ($\tau \sqcup \sigma$))
| *collection-binop-type IntersectionOp* (*Bag* τ) (*Set* σ) (*Set* ($\tau \sqcup \sigma$))
| *collection-binop-type IntersectionOp* (*Bag* τ) (*Bag* σ) (*Bag* ($\tau \sqcup \sigma$))

| *collection-binop-type SetMinusOp* (*Set* τ) (*Set* σ) (*Set* τ)
| *collection-binop-type SymmetricDifferenceOp* (*Set* τ) (*Set* σ) (*Set* ($\tau \sqcup \sigma$))

| *element-type* $\tau\ \varrho \Longrightarrow$ *update-element-type* τ ($\varrho \sqcup \sigma$) $\upsilon \Longrightarrow$
collection-binop-type IncludingOp $\tau\ \sigma\ \upsilon$
| *element-type* $\tau\ \varrho \Longrightarrow \sigma \leq \varrho \Longrightarrow$
collection-binop-type ExcludingOp $\tau\ \sigma\ \tau$

| $\sigma \leq \tau \Longrightarrow$
collection-binop-type AppendOp (*OrderedSet* τ) σ (*OrderedSet* τ)
| $\sigma \leq \tau \Longrightarrow$
collection-binop-type AppendOp (*Sequence* τ) σ (*Sequence* τ)
| $\sigma \leq \tau \Longrightarrow$
collection-binop-type PrependOp (*OrderedSet* τ) σ (*OrderedSet* τ)
| $\sigma \leq \tau \Longrightarrow$
collection-binop-type PrependOp (*Sequence* τ) σ (*Sequence* τ)

| $\sigma = \mathit{UnlimitedNatural}[1] - \mathit{Integer}[1] \Longrightarrow$
$\mathit{collection\text{-}binop\text{-}type}\ \mathit{CollectionAtOp}\ (\mathit{OrderedSet}\ \tau)\ \sigma\ \tau$
| $\sigma = \mathit{UnlimitedNatural}[1] - \mathit{Integer}[1] \Longrightarrow$
$\mathit{collection\text{-}binop\text{-}type}\ \mathit{CollectionAtOp}\ (\mathit{Sequence}\ \tau)\ \sigma\ \tau$

| $\sigma \leq \tau \Longrightarrow$
$\mathit{collection\text{-}binop\text{-}type}\ \mathit{CollectionIndexOfOp}\ (\mathit{OrderedSet}\ \tau)\ \sigma\ \mathit{Integer}[1]$
| $\sigma \leq \tau \Longrightarrow$
$\mathit{collection\text{-}binop\text{-}type}\ \mathit{CollectionIndexOfOp}\ (\mathit{Sequence}\ \tau)\ \sigma\ \mathit{Integer}[1]$

inductive *collection-ternop-type* **where**
$\sigma = \mathit{UnlimitedNatural}[1] - \mathit{Integer}[1] \Longrightarrow \varrho \leq \tau \Longrightarrow$
$\mathit{collection\text{-}ternop\text{-}type}\ \mathit{InsertAtOp}\ (\mathit{OrderedSet}\ \tau)\ \sigma\ \varrho\ (\mathit{OrderedSet}\ \tau)$
| $\sigma = \mathit{UnlimitedNatural}[1] - \mathit{Integer}[1] \Longrightarrow \varrho \leq \tau \Longrightarrow$
$\mathit{collection\text{-}ternop\text{-}type}\ \mathit{InsertAtOp}\ (\mathit{Sequence}\ \tau)\ \sigma\ \varrho\ (\mathit{Sequence}\ \tau)$
| $\sigma = \mathit{UnlimitedNatural}[1] - \mathit{Integer}[1] \Longrightarrow$
$\varrho = \mathit{UnlimitedNatural}[1] - \mathit{Integer}[1] \Longrightarrow$
$\mathit{collection\text{-}ternop\text{-}type}\ \mathit{SubOrderedSetOp}\ (\mathit{OrderedSet}\ \tau)\ \sigma\ \varrho\ (\mathit{OrderedSet}\ \tau)$
| $\sigma = \mathit{UnlimitedNatural}[1] - \mathit{Integer}[1] \Longrightarrow$
$\varrho = \mathit{UnlimitedNatural}[1] - \mathit{Integer}[1] \Longrightarrow$
$\mathit{collection\text{-}ternop\text{-}type}\ \mathit{SubSequenceOp}\ (\mathit{Sequence}\ \tau)\ \sigma\ \varrho\ (\mathit{Sequence}\ \tau)$

6.1.9 Coercions

inductive *unop-type* **where**
$\mathit{any\text{-}unop\text{-}type}\ \mathit{op}\ \tau\ \sigma \Longrightarrow$
$\mathit{unop\text{-}type}\ (\mathit{Inl}\ \mathit{op})\ \mathit{DotCall}\ \tau\ \sigma$
| $\mathit{boolean\text{-}unop\text{-}type}\ \mathit{op}\ \tau\ \sigma \Longrightarrow$
$\mathit{unop\text{-}type}\ (\mathit{Inr}\ (\mathit{Inl}\ \mathit{op}))\ \mathit{DotCall}\ \tau\ \sigma$
| $\mathit{numeric\text{-}unop\text{-}type}\ \mathit{op}\ \tau\ \sigma \Longrightarrow$
$\mathit{unop\text{-}type}\ (\mathit{Inr}\ (\mathit{Inr}\ (\mathit{Inl}\ \mathit{op})))\ \mathit{DotCall}\ \tau\ \sigma$
| $\mathit{string\text{-}unop\text{-}type}\ \mathit{op}\ \tau\ \sigma \Longrightarrow$
$\mathit{unop\text{-}type}\ (\mathit{Inr}\ (\mathit{Inr}\ (\mathit{Inr}\ (\mathit{Inl}\ \mathit{op}))))\ \mathit{DotCall}\ \tau\ \sigma$
| $\mathit{collection\text{-}unop\text{-}type}\ \mathit{op}\ \tau\ \sigma \Longrightarrow$
$\mathit{unop\text{-}type}\ (\mathit{Inr}\ (\mathit{Inr}\ (\mathit{Inr}\ (\mathit{Inr}\ \mathit{op}))))\ \mathit{ArrowCall}\ \tau\ \sigma$

inductive *binop-type* **where**
$\mathit{super\text{-}binop\text{-}type}\ \mathit{op}\ \tau\ \sigma\ \varrho \Longrightarrow$
$\mathit{binop\text{-}type}\ (\mathit{Inl}\ \mathit{op})\ \mathit{DotCall}\ \tau\ \sigma\ \varrho$
| $\mathit{boolean\text{-}binop\text{-}type}\ \mathit{op}\ \tau\ \sigma\ \varrho \Longrightarrow$
$\mathit{binop\text{-}type}\ (\mathit{Inr}\ (\mathit{Inl}\ \mathit{op}))\ \mathit{DotCall}\ \tau\ \sigma\ \varrho$
| $\mathit{numeric\text{-}binop\text{-}type}\ \mathit{op}\ \tau\ \sigma\ \varrho \Longrightarrow$
$\mathit{binop\text{-}type}\ (\mathit{Inr}\ (\mathit{Inr}\ (\mathit{Inl}\ \mathit{op})))\ \mathit{DotCall}\ \tau\ \sigma\ \varrho$
| $\mathit{string\text{-}binop\text{-}type}\ \mathit{op}\ \tau\ \sigma\ \varrho \Longrightarrow$
$\mathit{binop\text{-}type}\ (\mathit{Inr}\ (\mathit{Inr}\ (\mathit{Inr}\ (\mathit{Inl}\ \mathit{op}))))\ \mathit{DotCall}\ \tau\ \sigma\ \varrho$
| $\mathit{collection\text{-}binop\text{-}type}\ \mathit{op}\ \tau\ \sigma\ \varrho \Longrightarrow$
$\mathit{binop\text{-}type}\ (\mathit{Inr}\ (\mathit{Inr}\ (\mathit{Inr}\ (\mathit{Inr}\ \mathit{op}))))\ \mathit{ArrowCall}\ \tau\ \sigma\ \varrho$

inductive *ternop-type* **where**
string-ternop-type op τ σ ϱ υ ⟹
ternop-type (Inl op) DotCall τ σ ϱ υ
| *collection-ternop-type op τ σ ϱ υ ⟹*
ternop-type (Inr op) ArrowCall τ σ ϱ υ

inductive *op-type* **where**
unop-type op k τ υ ⟹
op-type (Inl op) k τ [] υ
| *binop-type op k τ σ υ ⟹*
op-type (Inr (Inl op)) k τ [σ] υ
| *ternop-type op k τ σ ϱ υ ⟹*
op-type (Inr (Inr (Inl op))) k τ [σ, ϱ] υ
| *operation τ op π oper ⟹*
op-type (Inr (Inr (Inr op))) DotCall τ π (oper-type oper)

6.1.10 Simplification Rules

inductive-simps *op-type-alt-simps*:
mataop-type τ op σ
typeop-type k op τ σ ϱ

op-type op k τ π σ
unop-type op k τ σ
binop-type op k τ σ ϱ
ternop-type op k τ σ ϱ υ

any-unop-type op τ σ
boolean-unop-type op τ σ
numeric-unop-type op τ σ
string-unop-type op τ σ
collection-unop-type op τ σ

super-binop-type op τ σ ϱ
boolean-binop-type op τ σ ϱ
numeric-binop-type op τ σ ϱ
string-binop-type op τ σ ϱ
collection-binop-type op τ σ ϱ

string-ternop-type op τ σ ϱ υ
collection-ternop-type op τ σ ϱ υ

6.1.11 Determinism

lemma *typeop-type-det*:
typeop-type op k τ σ ϱ_1 ⟹
typeop-type op k τ σ ϱ_2 ⟹ $\varrho_1 = \varrho_2$
by (*induct rule*: *typeop-type.induct*;
auto simp add: *typeop-type.simps update-element-type-det*)

lemma *any-unop-type-det*:
any-unop-type op τ $\sigma_1 \Longrightarrow$
any-unop-type op τ $\sigma_2 \Longrightarrow \sigma_1 = \sigma_2$
by (*induct rule*: *any-unop-type.induct*; *simp add*: *any-unop-type.simps*)

lemma *boolean-unop-type-det*:
boolean-unop-type op τ $\sigma_1 \Longrightarrow$
boolean-unop-type op τ $\sigma_2 \Longrightarrow \sigma_1 = \sigma_2$
by (*induct rule*: *boolean-unop-type.induct*; *simp add*: *boolean-unop-type.simps*)

lemma *numeric-unop-type-det*:
numeric-unop-type op τ $\sigma_1 \Longrightarrow$
numeric-unop-type op τ $\sigma_2 \Longrightarrow \sigma_1 = \sigma_2$
by (*induct rule*: *numeric-unop-type.induct*; *auto simp add*: *numeric-unop-type.simps*)

lemma *string-unop-type-det*:
string-unop-type op τ $\sigma_1 \Longrightarrow$
string-unop-type op τ $\sigma_2 \Longrightarrow \sigma_1 = \sigma_2$
by (*induct rule*: *string-unop-type.induct*; *simp add*: *string-unop-type.simps*)

lemma *collection-unop-type-det*:
collection-unop-type op τ $\sigma_1 \Longrightarrow$
collection-unop-type op τ $\sigma_2 \Longrightarrow \sigma_1 = \sigma_2$
apply (*induct rule*: *collection-unop-type.induct*)
by (*erule collection-unop-type.cases*;
auto simp add: *element-type-det update-element-type-det*)+

lemma *unop-type-det*:
unop-type op k τ $\sigma_1 \Longrightarrow$
unop-type op k τ $\sigma_2 \Longrightarrow \sigma_1 = \sigma_2$
by (*induct rule*: *unop-type.induct*;
simp add: *unop-type.simps any-unop-type-det*
boolean-unop-type-det numeric-unop-type-det
string-unop-type-det collection-unop-type-det)

lemma *super-binop-type-det*:
super-binop-type op τ σ $\varrho_1 \Longrightarrow$
super-binop-type op τ σ $\varrho_2 \Longrightarrow \varrho_1 = \varrho_2$
by (*induct rule*: *super-binop-type.induct*; *auto simp add*: *super-binop-type.simps*)

lemma *boolean-binop-type-det*:
boolean-binop-type op τ σ $\varrho_1 \Longrightarrow$
boolean-binop-type op τ σ $\varrho_2 \Longrightarrow \varrho_1 = \varrho_2$
by (*induct rule*: *boolean-binop-type.induct*; *simp add*: *boolean-binop-type.simps*)

lemma *numeric-binop-type-det*:
numeric-binop-type op τ σ $\varrho_1 \Longrightarrow$
numeric-binop-type op τ σ $\varrho_2 \Longrightarrow \varrho_1 = \varrho_2$
by (*induct rule*: *numeric-binop-type.induct*;

auto simp add: *numeric-binop-type.simps split*: *if-splits*)

lemma *string-binop-type-det*:
string-binop-type op τ σ ϱ_1 $\Longrightarrow$
string-binop-type op τ σ ϱ_2 $\Longrightarrow$ $\varrho_1 = \varrho_2$
by (*induct rule*: *string-binop-type.induct*; *simp add*: *string-binop-type.simps*)

lemma *collection-binop-type-det*:
collection-binop-type op τ σ ϱ_1 $\Longrightarrow$
collection-binop-type op τ σ ϱ_2 $\Longrightarrow$ $\varrho_1 = \varrho_2$
apply (*induct rule*: *collection-binop-type.induct*; *simp add*: *collection-binop-type.simps*)
using *element-type-det update-element-type-det* **by** *blast+*

lemma *binop-type-det*:
binop-type op k τ σ ϱ_1 $\Longrightarrow$
binop-type op k τ σ ϱ_2 $\Longrightarrow$ $\varrho_1 = \varrho_2$
by (*induct rule*: *binop-type.induct*;
simp add: *binop-type.simps super-binop-type-det*
boolean-binop-type-det numeric-binop-type-det
string-binop-type-det collection-binop-type-det)

lemma *string-ternop-type-det*:
string-ternop-type op τ σ ϱ υ_1 $\Longrightarrow$
string-ternop-type op τ σ ϱ υ_2 $\Longrightarrow$ $\upsilon_1 = \upsilon_2$
by (*induct rule*: *string-ternop-type.induct*; *simp add*: *string-ternop-type.simps*)

lemma *collection-ternop-type-det*:
collection-ternop-type op τ σ ϱ υ_1 $\Longrightarrow$
collection-ternop-type op τ σ ϱ υ_2 $\Longrightarrow$ $\upsilon_1 = \upsilon_2$
by (*induct rule*: *collection-ternop-type.induct*; *simp add*: *collection-ternop-type.simps*)

lemma *ternop-type-det*:
ternop-type op k τ σ ϱ υ_1 $\Longrightarrow$
ternop-type op k τ σ ϱ υ_2 $\Longrightarrow$ $\upsilon_1 = \upsilon_2$
by (*induct rule*: *ternop-type.induct*;
simp add: *ternop-type.simps string-ternop-type-det collection-ternop-type-det*)

lemma *op-type-det*:
op-type op k τ π σ $\Longrightarrow$
op-type op k τ π ϱ $\Longrightarrow$ $\sigma = \varrho$
apply (*induct rule*: *op-type.induct*)
apply (*erule op-type.cases*; *simp add*: *unop-type-det*)
apply (*erule op-type.cases*; *simp add*: *binop-type-det*)
apply (*erule op-type.cases*; *simp add*: *ternop-type-det*)
by (*erule op-type.cases*; *simp*; *metis operation-det*)

6.2 Expressions Typing

The following typing rules are preliminary. The final rules are given at the end of the next chapter.

inductive *typing* :: ('*a* :: *ocl-object-model*) *type env* ⇒ '*a expr* ⇒ '*a type* ⇒ *bool*
(‹(*1-*/ ⊢$_E$/ (- :/ -))› [*51*,*51*,*51*] *50*)
and *collection-parts-typing* (‹(*1-*/ ⊢$_C$/ (- :/ -))› [*51*,*51*,*51*] *50*)
and *collection-part-typing* (‹(*1-*/ ⊢$_P$/ (- :/ -))› [*51*,*51*,*51*] *50*)
and *iterator-typing* (‹(*1-*/ ⊢$_I$/ (- :/ -))› [*51*,*51*,*51*] *50*)
and *expr-list-typing* (‹(*1-*/ ⊢$_L$/ (- :/ -))› [*51*,*51*,*51*] *50*) **where**

— Primitive Literals

NullLiteralT:
Γ ⊢$_E$ *NullLiteral* : *OclVoid*[*?*]
|*BooleanLiteralT*:
Γ ⊢$_E$ *BooleanLiteral c* : *Boolean*[*1*]
|*RealLiteralT*:
Γ ⊢$_E$ *RealLiteral c* : *Real*[*1*]
|*IntegerLiteralT*:
Γ ⊢$_E$ *IntegerLiteral c* : *Integer*[*1*]
|*UnlimitedNaturalLiteralT*:
Γ ⊢$_E$ *UnlimitedNaturalLiteral c* : *UnlimitedNatural*[*1*]
|*StringLiteralT*:
Γ ⊢$_E$ *StringLiteral c* : *String*[*1*]
|*EnumLiteralT*:
has-literal enum lit ⟹
Γ ⊢$_E$ *EnumLiteral enum lit* : (*Enum enum*)[*1*]

— Collection Literals

|*SetLiteralT*:
Γ ⊢$_C$ *prts* : τ ⟹
Γ ⊢$_E$ *CollectionLiteral SetKind prts* : *Set* τ
|*OrderedSetLiteralT*:
Γ ⊢$_C$ *prts* : τ ⟹
Γ ⊢$_E$ *CollectionLiteral OrderedSetKind prts* : *OrderedSet* τ
|*BagLiteralT*:
Γ ⊢$_C$ *prts* : τ ⟹
Γ ⊢$_E$ *CollectionLiteral BagKind prts* : *Bag* τ
|*SequenceLiteralT*:
Γ ⊢$_C$ *prts* : τ ⟹
Γ ⊢$_E$ *CollectionLiteral SequenceKind prts* : *Sequence* τ

— We prohibit empty collection literals, because their type is unclear. We could use *OclVoid*[*1*] element type for empty collections, but the typing rules will give wrong types for nested collections, because, for example, *OclVoid*[*1*] ⊔ *Set*(*Integer*[*1*]) = *OclSuper*

$|CollectionPartsSingletonT$:
$\Gamma \vdash_P x : \tau \Longrightarrow$
$\Gamma \vdash_C [x] : \tau$
$|CollectionPartsListT$:
$\Gamma \vdash_P x : \tau \Longrightarrow$
$\Gamma \vdash_C y \# xs : \sigma \Longrightarrow$
$\Gamma \vdash_C x \# y \# xs : \tau \sqcup \sigma$

$|CollectionPartItemT$:
$\Gamma \vdash_E a : \tau \Longrightarrow$
$\Gamma \vdash_P CollectionItem\ a : \tau$
$|CollectionPartRangeT$:
$\Gamma \vdash_E a : \tau \Longrightarrow$
$\Gamma \vdash_E b : \sigma \Longrightarrow$
$\tau = UnlimitedNatural[1] - Integer[1] \Longrightarrow$
$\sigma = UnlimitedNatural[1] - Integer[1] \Longrightarrow$
$\Gamma \vdash_P CollectionRange\ a\ b : Integer[1]$

— Tuple Literals
— We do not prohibit empty tuples, because it could be useful. *Tuple*() is a supertype of all other tuple types.

$|TupleLiteralNilT$:
$\Gamma \vdash_E TupleLiteral\ [] : Tuple\ fmempty$
$|TupleLiteralConsT$:
$\Gamma \vdash_E TupleLiteral\ elems : Tuple\ \xi \Longrightarrow$
$\Gamma \vdash_E tuple\text{-}element\text{-}expr\ el : \tau \Longrightarrow$
$tuple\text{-}element\text{-}type\ el = Some\ \sigma \Longrightarrow$
$\tau \leq \sigma \Longrightarrow$
$\Gamma \vdash_E TupleLiteral\ (el\ \#\ elems) : Tuple\ (\xi(tuple\text{-}element\text{-}name\ el \mapsto_f \sigma))$

— Misc Expressions

$|LetT$:
$\Gamma \vdash_E init : \sigma \Longrightarrow$
$\sigma \leq \tau \Longrightarrow$
$\Gamma(v \mapsto_f \tau) \vdash_E body : \varrho \Longrightarrow$
$\Gamma \vdash_E Let\ v\ (Some\ \tau)\ init\ body : \varrho$
$|VarT$:
$fmlookup\ \Gamma\ v = Some\ \tau \Longrightarrow$
$\Gamma \vdash_E Var\ v : \tau$
$|IfT$:
$\Gamma \vdash_E a : Boolean[1] \Longrightarrow$
$\Gamma \vdash_E b : \sigma \Longrightarrow$
$\Gamma \vdash_E c : \varrho \Longrightarrow$
$\Gamma \vdash_E If\ a\ b\ c : \sigma \sqcup \varrho$

— Call Expressions

$|MetaOperationCallT$:
$mataop\text{-}type\ \tau\ op\ \sigma \Longrightarrow$
$\Gamma \vdash_E MetaOperationCall\ \tau\ op : \sigma$
$|StaticOperationCallT$:
$\Gamma \vdash_L params : \pi \Longrightarrow$
$static\text{-}operation\ \tau\ op\ \pi\ oper \Longrightarrow$
$\Gamma \vdash_E StaticOperationCall\ \tau\ op\ params : oper\text{-}type\ oper$

$|TypeOperationCallT$:
$\Gamma \vdash_E a : \tau \Longrightarrow$
$typeop\text{-}type\ k\ op\ \tau\ \sigma\ \varrho \Longrightarrow$
$\Gamma \vdash_E TypeOperationCall\ a\ k\ op\ \sigma : \varrho$

$|AttributeCallT$:
$\Gamma \vdash_E src : \langle \mathcal{C} \rangle_{\mathcal{T}}[1] \Longrightarrow$
$attribute\ \mathcal{C}\ prop\ \mathcal{D}\ \tau \Longrightarrow$
$\Gamma \vdash_E AttributeCall\ src\ DotCall\ prop : \tau$
$|AssociationEndCallT$:
$\Gamma \vdash_E src : \langle \mathcal{C} \rangle_{\mathcal{T}}[1] \Longrightarrow$
$association\text{-}end\ \mathcal{C}\ from\ role\ \mathcal{D}\ end \Longrightarrow$
$\Gamma \vdash_E AssociationEndCall\ src\ DotCall\ from\ role : assoc\text{-}end\text{-}type\ end$
$|AssociationClassCallT$:
$\Gamma \vdash_E src : \langle \mathcal{C} \rangle_{\mathcal{T}}[1] \Longrightarrow$
$referred\text{-}by\text{-}association\text{-}class\ \mathcal{C}\ from\ \mathcal{A}\ \mathcal{D} \Longrightarrow$
$\Gamma \vdash_E AssociationClassCall\ src\ DotCall\ from\ \mathcal{A} : class\text{-}assoc\text{-}type\ \mathcal{A}$
$|AssociationClassEndCallT$:
$\Gamma \vdash_E src : \langle \mathcal{A} \rangle_{\mathcal{T}}[1] \Longrightarrow$
$association\text{-}class\text{-}end\ \mathcal{A}\ role\ end \Longrightarrow$
$\Gamma \vdash_E AssociationClassEndCall\ src\ DotCall\ role : class\text{-}assoc\text{-}end\text{-}type\ end$
$|OperationCallT$:
$\Gamma \vdash_E src : \tau \Longrightarrow$
$\Gamma \vdash_L params : \pi \Longrightarrow$
$op\text{-}type\ op\ k\ \tau\ \pi\ \sigma \Longrightarrow$
$\Gamma \vdash_E OperationCall\ src\ k\ op\ params : \sigma$

$|TupleElementCallT$:
$\Gamma \vdash_E src : Tuple\ \pi \Longrightarrow$
$fmlookup\ \pi\ elem = Some\ \tau \Longrightarrow$
$\Gamma \vdash_E TupleElementCall\ src\ DotCall\ elem : \tau$

— Iterator Expressions

$|IteratorT$:
$\Gamma \vdash_E src : \tau \Longrightarrow$
$element\text{-}type\ \tau\ \sigma \Longrightarrow$
$\sigma \leq its\text{-}ty \Longrightarrow$
$\Gamma +\!\!+_f fmap\text{-}of\text{-}list\ (map\ (\lambda it.\ (it,\ its\text{-}ty))\ its) \vdash_E body : \varrho \Longrightarrow$
$\Gamma \vdash_I (src,\ its,\ (Some\ its\text{-}ty),\ body) : (\tau,\ \sigma,\ \varrho)$

$|IterateT$:
$\Gamma \vdash_I (src, its, its\text{-}ty, Let\ res\ (Some\ res\text{-}t)\ res\text{-}init\ body) : (\tau, \sigma, \varrho) \Longrightarrow$
$\varrho \leq res\text{-}t \Longrightarrow$
$\Gamma \vdash_E IterateCall\ src\ ArrowCall\ its\ its\text{-}ty\ res\ (Some\ res\text{-}t)\ res\text{-}init\ body : \varrho$

$|AnyIteratorT$:
$\Gamma \vdash_I (src, its, its\text{-}ty, body) : (\tau, \sigma, \varrho) \Longrightarrow$
$length\ its \leq 1 \Longrightarrow$
$\varrho \leq Boolean[?] \Longrightarrow$
$\Gamma \vdash_E AnyIteratorCall\ src\ ArrowCall\ its\ its\text{-}ty\ body : \sigma$
$|ClosureIteratorT$:
$\Gamma \vdash_I (src, its, its\text{-}ty, body) : (\tau, \sigma, \varrho) \Longrightarrow$
$length\ its \leq 1 \Longrightarrow$
$to\text{-}single\text{-}type\ \varrho \leq \sigma \Longrightarrow$
$to\text{-}unique\text{-}collection\ \tau\ \upsilon \Longrightarrow$
$\Gamma \vdash_E ClosureIteratorCall\ src\ ArrowCall\ its\ its\text{-}ty\ body : \upsilon$
$|CollectIteratorT$:
$\Gamma \vdash_I (src, its, its\text{-}ty, body) : (\tau, \sigma, \varrho) \Longrightarrow$
$length\ its \leq 1 \Longrightarrow$
$to\text{-}nonunique\text{-}collection\ \tau\ \upsilon \Longrightarrow$
$update\text{-}element\text{-}type\ \upsilon\ (to\text{-}single\text{-}type\ \varrho)\ \varphi \Longrightarrow$
$\Gamma \vdash_E CollectIteratorCall\ src\ ArrowCall\ its\ its\text{-}ty\ body : \varphi$
$|CollectNestedIteratorT$:
$\Gamma \vdash_I (src, its, its\text{-}ty, body) : (\tau, \sigma, \varrho) \Longrightarrow$
$length\ its \leq 1 \Longrightarrow$
$to\text{-}nonunique\text{-}collection\ \tau\ \upsilon \Longrightarrow$
$update\text{-}element\text{-}type\ \upsilon\ \varrho\ \varphi \Longrightarrow$
$\Gamma \vdash_E CollectNestedIteratorCall\ src\ ArrowCall\ its\ its\text{-}ty\ body : \varphi$
$|ExistsIteratorT$:
$\Gamma \vdash_I (src, its, its\text{-}ty, body) : (\tau, \sigma, \varrho) \Longrightarrow$
$\varrho \leq Boolean[?] \Longrightarrow$
$\Gamma \vdash_E ExistsIteratorCall\ src\ ArrowCall\ its\ its\text{-}ty\ body : \varrho$
$|ForAllIteratorT$:
$\Gamma \vdash_I (src, its, its\text{-}ty, body) : (\tau, \sigma, \varrho) \Longrightarrow$
$\varrho \leq Boolean[?] \Longrightarrow$
$\Gamma \vdash_E ForAllIteratorCall\ src\ ArrowCall\ its\ its\text{-}ty\ body : \varrho$
$|OneIteratorT$:
$\Gamma \vdash_I (src, its, its\text{-}ty, body) : (\tau, \sigma, \varrho) \Longrightarrow$
$length\ its \leq 1 \Longrightarrow$
$\varrho \leq Boolean[?] \Longrightarrow$
$\Gamma \vdash_E OneIteratorCall\ src\ ArrowCall\ its\ its\text{-}ty\ body : Boolean[1]$
$|IsUniqueIteratorT$:
$\Gamma \vdash_I (src, its, its\text{-}ty, body) : (\tau, \sigma, \varrho) \Longrightarrow$
$length\ its \leq 1 \Longrightarrow$
$\Gamma \vdash_E IsUniqueIteratorCall\ src\ ArrowCall\ its\ its\text{-}ty\ body : Boolean[1]$
$|SelectIteratorT$:
$\Gamma \vdash_I (src, its, its\text{-}ty, body) : (\tau, \sigma, \varrho) \Longrightarrow$
$length\ its \leq 1 \Longrightarrow$
$\varrho \leq Boolean[?] \Longrightarrow$

$\Gamma \vdash_E$ *SelectIteratorCall src ArrowCall its its-ty body* : τ
|*RejectIteratorT*:
$\Gamma \vdash_I$ (*src*, *its*, *its-ty*, *body*) : (τ, σ, ϱ) $\Longrightarrow$
length its $\leq$ *1* $\Longrightarrow$
$\varrho \leq$ *Boolean*[*?*] $\Longrightarrow$
$\Gamma \vdash_E$ *RejectIteratorCall src ArrowCall its its-ty body* : τ
|*SortedByIteratorT*:
$\Gamma \vdash_I$ (*src*, *its*, *its-ty*, *body*) : (τ, σ, ϱ) $\Longrightarrow$
length its $\leq$ *1* $\Longrightarrow$
to-ordered-collection τ υ $\Longrightarrow$
$\Gamma \vdash_E$ *SortedByIteratorCall src ArrowCall its its-ty body* : υ

— Expression Lists

|*ExprListNilT*:
$\Gamma \vdash_L$ [] : []
|*ExprListConsT*:
$\Gamma \vdash_E$ *expr* : τ $\Longrightarrow$
$\Gamma \vdash_L$ *exprs* : π $\Longrightarrow$
$\Gamma \vdash_L$ *expr* # *exprs* : τ # π

6.3 Elimination Rules

inductive-cases *NullLiteralTE* [*elim*]: $\Gamma \vdash_E$ *NullLiteral* : τ
inductive-cases *BooleanLiteralTE* [*elim*]: $\Gamma \vdash_E$ *BooleanLiteral c* : τ
inductive-cases *RealLiteralTE* [*elim*]: $\Gamma \vdash_E$ *RealLiteral c* : τ
inductive-cases *IntegerLiteralTE* [*elim*]: $\Gamma \vdash_E$ *IntegerLiteral c* : τ
inductive-cases *UnlimitedNaturalLiteralTE* [*elim*]: $\Gamma \vdash_E$ *UnlimitedNaturalLiteral c* : τ
inductive-cases *StringLiteralTE* [*elim*]: $\Gamma \vdash_E$ *StringLiteral c* : τ
inductive-cases *EnumLiteralTE* [*elim*]: $\Gamma \vdash_E$ *EnumLiteral enm lit* : τ
inductive-cases *CollectionLiteralTE* [*elim*]: $\Gamma \vdash_E$ *CollectionLiteral k prts* : τ
inductive-cases *TupleLiteralTE* [*elim*]: $\Gamma \vdash_E$ *TupleLiteral elems* : τ

inductive-cases *LetTE* [*elim*]: $\Gamma \vdash_E$ *Let v* τ *init body* : σ
inductive-cases *VarTE* [*elim*]: $\Gamma \vdash_E$ *Var v* : τ
inductive-cases *IfTE* [*elim*]: $\Gamma \vdash_E$ *If a b c* : τ

inductive-cases *MetaOperationCallTE* [*elim*]: $\Gamma \vdash_E$ *MetaOperationCall* τ *op* : σ

inductive-cases *StaticOperationCallTE* [*elim*]: $\Gamma \vdash_E$ *StaticOperationCall* τ *op as* : σ

inductive-cases *TypeOperationCallTE* [*elim*]: $\Gamma \vdash_E$ *TypeOperationCall a k op* σ : τ
inductive-cases *AttributeCallTE* [*elim*]: $\Gamma \vdash_E$ *AttributeCall src k prop* : τ
inductive-cases *AssociationEndCallTE* [*elim*]: $\Gamma \vdash_E$ *AssociationEndCall src k role from* : τ
inductive-cases *AssociationClassCallTE* [*elim*]: $\Gamma \vdash_E$ *AssociationClassCall src k*

a from : τ
inductive-cases *AssociationClassEndCallTE* [*elim*]: $\Gamma \vdash_E$ *AssociationClassEndCall src k role* : τ
inductive-cases *OperationCallTE* [*elim*]: $\Gamma \vdash_E$ *OperationCall src k op params* : τ
inductive-cases *TupleElementCallTE* [*elim*]: $\Gamma \vdash_E$ *TupleElementCall src k elem* : τ

inductive-cases *IteratorTE* [*elim*]: $\Gamma \vdash_I$ (*src*, *its*, *body*) : *ys*
inductive-cases *IterateTE* [*elim*]: $\Gamma \vdash_E$ *IterateCall src k its its-ty res res-t res-init body* : τ
inductive-cases *AnyIteratorTE* [*elim*]: $\Gamma \vdash_E$ *AnyIteratorCall src k its its-ty body* : τ
inductive-cases *ClosureIteratorTE* [*elim*]: $\Gamma \vdash_E$ *ClosureIteratorCall src k its its-ty body* : τ
inductive-cases *CollectIteratorTE* [*elim*]: $\Gamma \vdash_E$ *CollectIteratorCall src k its its-ty body* : τ
inductive-cases *CollectNestedIteratorTE* [*elim*]: $\Gamma \vdash_E$ *CollectNestedIteratorCall src k its its-ty body* : τ
inductive-cases *ExistsIteratorTE* [*elim*]: $\Gamma \vdash_E$ *ExistsIteratorCall src k its its-ty body* : τ
inductive-cases *ForAllIteratorTE* [*elim*]: $\Gamma \vdash_E$ *ForAllIteratorCall src k its its-ty body* : τ
inductive-cases *OneIteratorTE* [*elim*]: $\Gamma \vdash_E$ *OneIteratorCall src k its its-ty body* : τ
inductive-cases *IsUniqueIteratorTE* [*elim*]: $\Gamma \vdash_E$ *IsUniqueIteratorCall src k its its-ty body* : τ
inductive-cases *SelectIteratorTE* [*elim*]: $\Gamma \vdash_E$ *SelectIteratorCall src k its its-ty body* : τ
inductive-cases *RejectIteratorTE* [*elim*]: $\Gamma \vdash_E$ *RejectIteratorCall src k its its-ty body* : τ
inductive-cases *SortedByIteratorTE* [*elim*]: $\Gamma \vdash_E$ *SortedByIteratorCall src k its its-ty body* : τ

inductive-cases *CollectionPartsNilTE* [*elim*]: $\Gamma \vdash_C [x] : \tau$
inductive-cases *CollectionPartsItemTE* [*elim*]: $\Gamma \vdash_C x \# y \# xs : \tau$

inductive-cases *CollectionItemTE* [*elim*]: $\Gamma \vdash_P$ *CollectionItem a* : τ
inductive-cases *CollectionRangeTE* [*elim*]: $\Gamma \vdash_P$ *CollectionRange a b* : τ

inductive-cases *ExprListTE* [*elim*]: $\Gamma \vdash_L$ *exprs* : π

6.4 Simplification Rules

inductive-simps *typing-alt-simps*:
$\Gamma \vdash_E$ *NullLiteral* : τ
$\Gamma \vdash_E$ *BooleanLiteral c* : τ
$\Gamma \vdash_E$ *RealLiteral c* : τ
$\Gamma \vdash_E$ *UnlimitedNaturalLiteral c* : τ

$\Gamma \vdash_E \mathit{IntegerLiteral}\ c : \tau$
$\Gamma \vdash_E \mathit{StringLiteral}\ c : \tau$
$\Gamma \vdash_E \mathit{EnumLiteral}\ enm\ lit : \tau$
$\Gamma \vdash_E \mathit{CollectionLiteral}\ k\ prts : \tau$
$\Gamma \vdash_E \mathit{TupleLiteral}\ [] : \tau$
$\Gamma \vdash_E \mathit{TupleLiteral}\ (x \# xs) : \tau$

$\Gamma \vdash_E \mathit{Let}\ v\ \tau\ init\ body : \sigma$
$\Gamma \vdash_E \mathit{Var}\ v : \tau$
$\Gamma \vdash_E \mathit{If}\ a\ b\ c : \tau$

$\Gamma \vdash_E \mathit{MetaOperationCall}\ \tau\ op : \sigma$
$\Gamma \vdash_E \mathit{StaticOperationCall}\ \tau\ op\ as : \sigma$

$\Gamma \vdash_E \mathit{TypeOperationCall}\ a\ k\ op\ \sigma : \tau$
$\Gamma \vdash_E \mathit{AttributeCall}\ src\ k\ prop : \tau$
$\Gamma \vdash_E \mathit{AssociationEndCall}\ src\ k\ role\ from : \tau$
$\Gamma \vdash_E \mathit{AssociationClassCall}\ src\ k\ a\ from : \tau$
$\Gamma \vdash_E \mathit{AssociationClassEndCall}\ src\ k\ role : \tau$
$\Gamma \vdash_E \mathit{OperationCall}\ src\ k\ op\ params : \tau$
$\Gamma \vdash_E \mathit{TupleElementCall}\ src\ k\ elem : \tau$

$\Gamma \vdash_I (src, its, body) : ys$
$\Gamma \vdash_E \mathit{IterateCall}\ src\ k\ its\ its\text{-}ty\ res\ res\text{-}t\ res\text{-}init\ body : \tau$
$\Gamma \vdash_E \mathit{AnyIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$
$\Gamma \vdash_E \mathit{ClosureIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$
$\Gamma \vdash_E \mathit{CollectIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$
$\Gamma \vdash_E \mathit{CollectNestedIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$
$\Gamma \vdash_E \mathit{ExistsIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$
$\Gamma \vdash_E \mathit{ForAllIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$
$\Gamma \vdash_E \mathit{OneIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$
$\Gamma \vdash_E \mathit{IsUniqueIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$
$\Gamma \vdash_E \mathit{SelectIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$
$\Gamma \vdash_E \mathit{RejectIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$
$\Gamma \vdash_E \mathit{SortedByIteratorCall}\ src\ k\ its\ its\text{-}ty\ body : \tau$

$\Gamma \vdash_C [x] : \tau$
$\Gamma \vdash_C x \# y \# xs : \tau$

$\Gamma \vdash_P \mathit{CollectionItem}\ a : \tau$
$\Gamma \vdash_P \mathit{CollectionRange}\ a\ b : \tau$

$\Gamma \vdash_L [] : \pi$
$\Gamma \vdash_L x \# xs : \pi$

6.5 Determinism

lemma
typing-det:

```
    Γ ⊢E expr : τ ⟹
    Γ ⊢E expr : σ ⟹ τ = σ  and
  collection-parts-typing-det:
    Γ ⊢C prts : τ ⟹
    Γ ⊢C prts : σ ⟹ τ = σ  and
  collection-part-typing-det:
    Γ ⊢P prt : τ ⟹
    Γ ⊢P prt : σ ⟹ τ = σ  and
  iterator-typing-det:
    Γ ⊢I (src, its, body) : xs ⟹
    Γ ⊢I (src, its, body) : ys ⟹ xs = ys  and
  expr-list-typing-det:
    Γ ⊢L exprs : π ⟹
    Γ ⊢L exprs : ξ ⟹ π = ξ
proof (induct arbitrary: σ and σ and σ and ys and ξ
    rule: typing-collection-parts-typing-collection-part-typing-iterator-typing-expr-list-typing.inducts)
  case (NullLiteralT Γ) thus ?case by auto
next
  case (BooleanLiteralT Γ c) thus ?case by auto
next
  case (RealLiteralT Γ c) thus ?case by auto
next
  case (IntegerLiteralT Γ c) thus ?case by auto
next
  case (UnlimitedNaturalLiteralT Γ c) thus ?case by auto
next
  case (StringLiteralT Γ c) thus ?case by auto
next
  case (EnumLiteralT Γ τ lit) thus ?case by auto
next
  case (SetLiteralT Γ prts τ) thus ?case by blast
next
  case (OrderedSetLiteralT Γ prts τ) thus ?case by blast
next
  case (BagLiteralT Γ prts τ) thus ?case by blast
next
  case (SequenceLiteralT Γ prts τ) thus ?case by blast
next
  case (CollectionPartsSingletonT Γ x τ) thus ?case by blast
next
  case (CollectionPartsListT Γ x τ y xs σ) thus ?case by blast
next
  case (CollectionPartItemT Γ a τ) thus ?case by blast
next
  case (CollectionPartRangeT Γ a τ b σ) thus ?case by blast
next
  case (TupleLiteralNilT Γ) thus ?case by auto
next
  case (TupleLiteralConsT Γ elems ξ el τ) show ?case
```

```
    apply (insert TupleLiteralConsT.prems)
    apply (erule TupleLiteralTE, simp)
    using TupleLiteralConsT.hyps by auto
next
  case (LetT Γ 𝓜 init σ τ v body ϱ) thus ?case by blast
next
  case (VarT Γ v τ 𝓜) thus ?case by auto
next
  case (IfT Γ a τ b σ c ϱ) thus ?case
    apply (insert IfT.prems)
    apply (erule IfTE)
    by (simp add: IfT.hyps)
next
  case (MetaOperationCallT τ op σ Γ) thus ?case
    by (metis MetaOperationCallTE mataop-type.cases)
next
  case (StaticOperationCallT τ op π oper Γ as) thus ?case
    apply (insert StaticOperationCallT.prems)
    apply (erule StaticOperationCallTE)
    using StaticOperationCallT.hyps static-operation-det by blast
next
  case (TypeOperationCallT Γ a τ op σ ϱ) thus ?case
    by (metis TypeOperationCallTE typeop-type-det)
next
  case (AttributeCallT Γ src τ 𝓒 prop 𝓓 σ) show ?case
    apply (insert AttributeCallT.prems)
    apply (erule AttributeCallTE)
    using AttributeCallT.hyps attribute-det by blast
next
  case (AssociationEndCallT Γ src 𝓒 from role 𝓓 end ) show ?case
    apply (insert AssociationEndCallT.prems)
    apply (erule AssociationEndCallTE)
    using AssociationEndCallT.hyps association-end-det by blast
next
  case (AssociationClassCallT Γ src 𝓒 from 𝓐) thus ?case by blast
next
  case (AssociationClassEndCallT Γ src τ 𝓐 role end ) show ?case
    apply (insert AssociationClassEndCallT.prems)
    apply (erule AssociationClassEndCallTE)
    using AssociationClassEndCallT.hyps association-class-end-det by blast
next
  case (OperationCallT Γ src τ params π op k) show ?case
    apply (insert OperationCallT.prems)
    apply (erule OperationCallTE)
    using OperationCallT.hyps op-type-det by blast
next
  case (TupleElementCallT Γ src π elem τ) thus ?case
    apply (insert TupleElementCallT.prems)
    apply (erule TupleElementCallTE)
```

```
    using TupleElementCallT.hyps by fastforce
next
  case (IteratorT Γ src τ σ its-ty its body ϱ) show ?case
    apply (insert IteratorT.prems)
    apply (erule IteratorTE)
    using IteratorT.hyps element-type-det by blast
next
  case (IterateT Γ src its its-ty res res-t res-init body τ σ ϱ) show ?case
    apply (insert IterateT.prems)
    using IterateT.hyps by blast
next
  case (AnyIteratorT Γ src its its-ty body τ σ ϱ) thus ?case
    by (meson AnyIteratorTE Pair-inject)
next
  case (ClosureIteratorT Γ src its its-ty body τ σ ϱ υ) show ?case
    apply (insert ClosureIteratorT.prems)
    apply (erule ClosureIteratorTE)
    using ClosureIteratorT.hyps to-unique-collection-det by blast
next
  case (CollectIteratorT Γ src its its-ty body τ σ ϱ υ) show ?case
    apply (insert CollectIteratorT.prems)
    apply (erule CollectIteratorTE)
    using CollectIteratorT.hyps to-nonunique-collection-det
      update-element-type-det Pair-inject by metis
next
  case (CollectNestedIteratorT Γ src its its-ty body τ σ ϱ υ) show ?case
    apply (insert CollectNestedIteratorT.prems)
    apply (erule CollectNestedIteratorTE)
    using CollectNestedIteratorT.hyps to-nonunique-collection-det
      update-element-type-det Pair-inject by metis
next
  case (ExistsIteratorT Γ src its its-ty body τ σ ϱ) show ?case
    apply (insert ExistsIteratorT.prems)
    apply (erule ExistsIteratorTE)
    using ExistsIteratorT.hyps Pair-inject by metis
next
  case (ForAllIteratorT Γ M src its its-ty body τ σ ϱ) show ?case
    apply (insert ForAllIteratorT.prems)
    apply (erule ForAllIteratorTE)
    using ForAllIteratorT.hyps Pair-inject by metis
next
  case (OneIteratorT Γ M src its its-ty body τ σ ϱ) show ?case
    apply (insert OneIteratorT.prems)
    apply (erule OneIteratorTE)
    by simp
next
  case (IsUniqueIteratorT Γ M src its its-ty body τ σ ϱ) show ?case
    apply (insert IsUniqueIteratorT.prems)
    apply (erule IsUniqueIteratorTE)
```

```
    by simp
next
  case (SelectIteratorT Γ M src its its-ty body τ σ ϱ) show ?case
    apply (insert SelectIteratorT.prems)
    apply (erule SelectIteratorTE)
    using SelectIteratorT.hyps by blast
next
  case (RejectIteratorT Γ M src its its-ty body τ σ ϱ) show ?case
    apply (insert RejectIteratorT.prems)
    apply (erule RejectIteratorTE)
    using RejectIteratorT.hyps by blast
next
  case (SortedByIteratorT Γ M src its its-ty body τ σ ϱ υ) show ?case
    apply (insert SortedByIteratorT.prems)
    apply (erule SortedByIteratorTE)
    using SortedByIteratorT.hyps to-ordered-collection-det by blast
next
  case (ExprListNilT Γ) thus ?case
    using expr-list-typing.cases by auto
next
  case (ExprListConsT Γ expr τ exprs π) show ?case
    apply (insert ExprListConsT.prems)
    apply (erule ExprListTE)
    by (simp-all add: ExprListConsT.hyps)
qed
```

6.6 Code Setup

```
code-pred op-type .

code-pred (modes:
    i ⇒ i ⇒ i ⇒ bool,
    i ⇒ i ⇒ o ⇒ bool) iterator-typing .

end
```

Chapter 7

Normalization

theory *OCL-Normalization*
imports *OCL-Typing*
begin

7.1 Normalization Rules

The following expression normalization rules includes two kinds of an abstract syntax tree transformations:

- determination of implicit types of variables, iterators, and tuple elements,
- unfolding of navigation shorthands and safe navigation operators, described in Table 7.1.

The following variables are used in the table:

- `x` is a non-nullable value,
- `n` is a nullable value,
- `xs` is a collection of non-nullable values,
- `ns` is a collection of nullable values.

Please take a note that name resolution of variables, types, attributes, and associations is out of scope of this section. It should be done on a previous phase during transformation of a concrete syntax tree to an abstract syntax tree.

fun *string-of-nat* :: *nat* ⇒ *string* **where**
string-of-nat n = (*if n* < *10 then* [*char-of* (*48* + *n*)]
else string-of-nat (*n div 10*) @ [*char-of* (*48* + (*n mod 10*))])

definition *new-vname* ≡ *String.implode* ∘ *string-of-nat* ∘ *fcard* ∘ *fmdom*

Table 7.1: Expression Normalization Rules

Orig. expr.	Normalized expression
`x.op()`	`x.op()`
`n.op()`	`n.op()`*
`x?.op()`	—
`n?.op()`	`if n <> null then n.oclAsType(T[1]).op() else null endif`**
`x->op()`	`x.oclAsSet()->op()`
`n->op()`	`n.oclAsSet()->op()`
`x?->op()`	—
`n?->op()`	—
`xs.op()`	`xs->collect(x | x.op())`
`ns.op()`	`ns->collect(n | n.op())`*
`xs?.op()`	—
`ns?.op()`	`ns->selectByKind(T[1])->collect(x | x.op())`
`xs->op()`	`xs->op()`
`ns->op()`	`ns->op()`
`xs?->op()`	—
`ns?->op()`	`ns->selectByKind(T[1])->op()`

* The resulting expression will be ill-typed if the operation is unsafe. An unsafe operation is an operation which is well-typed for a non-nullable source only.

** It would be a good idea to prohibit such a transformation for safe operations. A safe operation is an operation which is well-typed for a nullable source. However, it is hard to define safe operations formally considering operations overloading, complex relations between operation parameters types (please see the typing rules for the equality operator), and user-defined operations.

inductive *normalize*
:: ('*a* :: *ocl-object-model*) *type env* ⇒ '*a expr* ⇒ '*a expr* ⇒ *bool*
(‹- ⊢ - ⇛/ -› [*51*,*51*,*51*] *50*) **and**
normalize-call (‹- ⊢$_C$ - ⇛/ -› [*51*,*51*,*51*] *50*) **and**
normalize-expr-list (‹- ⊢$_L$ - ⇛/ -› [*51*,*51*,*51*] *50*)
where
LiteralN:
Γ ⊢ *Literal a* ⇛ *Literal a*
|*ExplicitlyTypedLetN*:
Γ ⊢ *init*$_1$ ⇛ *init*$_2$ ⟹
Γ(*v* ↦$_f$ τ) ⊢ *body*$_1$ ⇛ *body*$_2$ ⟹
Γ ⊢ *Let v* (*Some* τ) *init*$_1$ *body*$_1$ ⇛ *Let v* (*Some* τ) *init*$_2$ *body*$_2$
|*ImplicitlyTypedLetN*:
Γ ⊢ *init*$_1$ ⇛ *init*$_2$ ⟹
Γ ⊢$_E$ *init*$_2$: τ ⟹
Γ(*v* ↦$_f$ τ) ⊢ *body*$_1$ ⇛ *body*$_2$ ⟹

$\Gamma \vdash \mathit{Let}\ v\ \mathit{None}\ \mathit{init}_1\ \mathit{body}_1 \Rrightarrow \mathit{Let}\ v\ (\mathit{Some}\ \tau)\ \mathit{init}_2\ \mathit{body}_2$
|*VarN*:
$\Gamma \vdash \mathit{Var}\ v \Rrightarrow \mathit{Var}\ v$
|*IfN*:
$\Gamma \vdash a_1 \Rrightarrow a_2 \Longrightarrow$
$\Gamma \vdash b_1 \Rrightarrow b_2 \Longrightarrow$
$\Gamma \vdash c_1 \Rrightarrow c_2 \Longrightarrow$
$\Gamma \vdash \mathit{If}\ a_1\ b_1\ c_1 \Rrightarrow \mathit{If}\ a_2\ b_2\ c_2$

|*MetaOperationCallN*:
$\Gamma \vdash \mathit{MetaOperationCall}\ \tau\ \mathit{op} \Rrightarrow \mathit{MetaOperationCall}\ \tau\ \mathit{op}$
|*StaticOperationCallN*:
$\Gamma \vdash_L \mathit{params}_1 \Rrightarrow \mathit{params}_2 \Longrightarrow$
$\Gamma \vdash \mathit{StaticOperationCall}\ \tau\ \mathit{op}\ \mathit{params}_1 \Rrightarrow \mathit{StaticOperationCall}\ \tau\ \mathit{op}\ \mathit{params}_2$

|*OclAnyDotCallN*:
$\Gamma \vdash \mathit{src}_1 \Rrightarrow \mathit{src}_2 \Longrightarrow$
$\Gamma \vdash_E \mathit{src}_2 : \tau \Longrightarrow$
$\tau \leq \mathit{OclAny}[?] \vee \tau \leq \mathit{Tuple}\ \mathit{fmempty} \Longrightarrow$
$(\Gamma, \tau) \vdash_C \mathit{call}_1 \Rrightarrow \mathit{call}_2 \Longrightarrow$
$\Gamma \vdash \mathit{Call}\ \mathit{src}_1\ \mathit{DotCall}\ \mathit{call}_1 \Rrightarrow \mathit{Call}\ \mathit{src}_2\ \mathit{DotCall}\ \mathit{call}_2$
|*OclAnySafeDotCallN*:
$\Gamma \vdash \mathit{src}_1 \Rrightarrow \mathit{src}_2 \Longrightarrow$
$\Gamma \vdash_E \mathit{src}_2 : \tau \Longrightarrow$
$\mathit{OclVoid}[?] \leq \tau \Longrightarrow$
$(\Gamma, \mathit{to\text{-}required\text{-}type}\ \tau) \vdash_C \mathit{call}_1 \Rrightarrow \mathit{call}_2 \Longrightarrow$
$\mathit{src}_3 = \mathit{TypeOperationCall}\ \mathit{src}_2\ \mathit{DotCall}\ \mathit{OclAsTypeOp}\ (\mathit{to\text{-}required\text{-}type}\ \tau) \Longrightarrow$
$\Gamma \vdash \mathit{Call}\ \mathit{src}_1\ \mathit{SafeDotCall}\ \mathit{call}_1 \Rrightarrow$
$\quad \mathit{If}\ (\mathit{OperationCall}\ \mathit{src}_2\ \mathit{DotCall}\ \mathit{NotEqualOp}\ [\mathit{NullLiteral}])$
$\quad\quad (\mathit{Call}\ \mathit{src}_3\ \mathit{DotCall}\ \mathit{call}_2)$
$\quad\quad \mathit{NullLiteral}$
|*OclAnyArrowCallN*:
$\Gamma \vdash \mathit{src}_1 \Rrightarrow \mathit{src}_2 \Longrightarrow$
$\Gamma \vdash_E \mathit{src}_2 : \tau \Longrightarrow$
$\tau \leq \mathit{OclAny}[?] \vee \tau \leq \mathit{Tuple}\ \mathit{fmempty} \Longrightarrow$
$\mathit{src}_3 = \mathit{OperationCall}\ \mathit{src}_2\ \mathit{DotCall}\ \mathit{OclAsSetOp}\ [] \Longrightarrow$
$\Gamma \vdash_E \mathit{src}_3 : \sigma \Longrightarrow$
$(\Gamma, \sigma) \vdash_C \mathit{call}_1 \Rrightarrow \mathit{call}_2 \Longrightarrow$
$\Gamma \vdash \mathit{Call}\ \mathit{src}_1\ \mathit{ArrowCall}\ \mathit{call}_1 \Rrightarrow \mathit{Call}\ \mathit{src}_3\ \mathit{ArrowCall}\ \mathit{call}_2$

|*CollectionArrowCallN*:
$\Gamma \vdash \mathit{src}_1 \Rrightarrow \mathit{src}_2 \Longrightarrow$
$\Gamma \vdash_E \mathit{src}_2 : \tau \Longrightarrow$
$\mathit{element\text{-}type}\ \tau\ \text{-} \Longrightarrow$
$(\Gamma, \tau) \vdash_C \mathit{call}_1 \Rrightarrow \mathit{call}_2 \Longrightarrow$
$\Gamma \vdash \mathit{Call}\ \mathit{src}_1\ \mathit{ArrowCall}\ \mathit{call}_1 \Rrightarrow \mathit{Call}\ \mathit{src}_2\ \mathit{ArrowCall}\ \mathit{call}_2$
|*CollectionSafeArrowCallN*:
$\Gamma \vdash \mathit{src}_1 \Rrightarrow \mathit{src}_2 \Longrightarrow$
$\Gamma \vdash_E \mathit{src}_2 : \tau \Longrightarrow$

$element\text{-}type\ \tau\ \sigma \Longrightarrow$
$OclVoid[?] \leq \sigma \Longrightarrow$
$src_3 = TypeOperationCall\ src_2\ ArrowCall\ SelectByKindOp$
$(to\text{-}required\text{-}type\ \sigma) \Longrightarrow$
$\Gamma \vdash_E src_3 : \varrho \Longrightarrow$
$(\Gamma, \varrho) \vdash_C call_1 \Rrightarrow call_2 \Longrightarrow$
$\Gamma \vdash Call\ src_1\ SafeArrowCall\ call_1 \Rrightarrow Call\ src_3\ ArrowCall\ call_2$

|*CollectionDotCallN*:
$\Gamma \vdash src_1 \Rrightarrow src_2 \Longrightarrow$
$\Gamma \vdash_E src_2 : \tau \Longrightarrow$
$element\text{-}type\ \tau\ \sigma \Longrightarrow$
$(\Gamma, \sigma) \vdash_C call_1 \Rrightarrow call_2 \Longrightarrow$
$it = new\text{-}vname\ \Gamma \Longrightarrow$
$\Gamma \vdash Call\ src_1\ DotCall\ call_1 \Rrightarrow$
$CollectIteratorCall\ src_2\ ArrowCall\ [it]\ (Some\ \sigma)\ (Call\ (Var\ it)\ DotCall\ call_2)$

|*CollectionSafeDotCallN*:
$\Gamma \vdash src_1 \Rrightarrow src_2 \Longrightarrow$
$\Gamma \vdash_E src_2 : \tau \Longrightarrow$
$element\text{-}type\ \tau\ \sigma \Longrightarrow$
$OclVoid[?] \leq \sigma \Longrightarrow$
$\varrho = to\text{-}required\text{-}type\ \sigma \Longrightarrow$
$src_3 = TypeOperationCall\ src_2\ ArrowCall\ SelectByKindOp\ \varrho \Longrightarrow$
$(\Gamma, \varrho) \vdash_C call_1 \Rrightarrow call_2 \Longrightarrow$
$it = new\text{-}vname\ \Gamma \Longrightarrow$
$\Gamma \vdash Call\ src_1\ SafeDotCall\ call_1 \Rrightarrow$
$CollectIteratorCall\ src_3\ ArrowCall\ [it]\ (Some\ \varrho)\ (Call\ (Var\ it)\ DotCall\ call_2)$

|*TypeOperationN*:
$(\Gamma, \tau) \vdash_C TypeOperation\ op\ ty \Rrightarrow TypeOperation\ op\ ty$
|*AttributeN*:
$(\Gamma, \tau) \vdash_C Attribute\ attr \Rrightarrow Attribute\ attr$
|*AssociationEndN*:
$(\Gamma, \tau) \vdash_C AssociationEnd\ role\ from \Rrightarrow AssociationEnd\ role\ from$
|*AssociationClassN*:
$(\Gamma, \tau) \vdash_C AssociationClass\ \mathcal{A}\ from \Rrightarrow AssociationClass\ \mathcal{A}\ from$
|*AssociationClassEndN*:
$(\Gamma, \tau) \vdash_C AssociationClassEnd\ role \Rrightarrow AssociationClassEnd\ role$
|*OperationN*:
$\Gamma \vdash_L params_1 \Rrightarrow params_2 \Longrightarrow$
$(\Gamma, \tau) \vdash_C Operation\ op\ params_1 \Rrightarrow Operation\ op\ params_2$
|*TupleElementN*:
$(\Gamma, \tau) \vdash_C TupleElement\ elem \Rrightarrow TupleElement\ elem$

|*ExplicitlyTypedIterateN*:
$\Gamma \vdash res\text{-}init_1 \Rrightarrow res\text{-}init_2 \Longrightarrow$
$\Gamma +\!\!+_f fmap\text{-}of\text{-}list\ (map\ (\lambda it.\ (it, \sigma))\ its) \vdash$
$Let\ res\ res\text{-}t_1\ res\text{-}init_1\ body_1 \Rrightarrow Let\ res\ res\text{-}t_2\ res\text{-}init_2\ body_2 \Longrightarrow$
$(\Gamma, \tau) \vdash_C Iterate\ its\ (Some\ \sigma)\ res\ res\text{-}t_1\ res\text{-}init_1\ body_1 \Rrightarrow$
$Iterate\ its\ (Some\ \sigma)\ res\ res\text{-}t_2\ res\text{-}init_2\ body_2$

|*ImplicitlyTypedIterateN*:
element-type τ σ $\Longrightarrow$
$\Gamma \vdash$ *res-init*$_1$ $\Rrightarrow$ *res-init*$_2$ $\Longrightarrow$
Γ $+\!\!+_f$ *fmap-of-list* (*map* (λ*it*. (*it*, σ)) *its*) $\vdash$
Let res res-t$_1$ *res-init*$_1$ *body*$_1$ $\Rrightarrow$ *Let res res-t*$_2$ *res-init*$_2$ *body*$_2$ $\Longrightarrow$
(Γ, τ) $\vdash_C$ *Iterate its None res res-t*$_1$ *res-init*$_1$ *body*$_1$ $\Rrightarrow$
Iterate its (*Some* σ) *res res-t*$_2$ *res-init*$_2$ *body*$_2$

|*ExplicitlyTypedIteratorN*:
Γ $+\!\!+_f$ *fmap-of-list* (*map* (λ*it*. (*it*, σ)) *its*) $\vdash$ *body*$_1$ $\Rrightarrow$ *body*$_2$ $\Longrightarrow$
(Γ, τ) $\vdash_C$ *Iterator iter its* (*Some* σ) *body*$_1$ $\Rrightarrow$ *Iterator iter its* (*Some* σ) *body*$_2$
|*ImplicitlyTypedIteratorN*:
element-type τ σ $\Longrightarrow$
Γ $+\!\!+_f$ *fmap-of-list* (*map* (λ*it*. (*it*, σ)) *its*) $\vdash$ *body*$_1$ $\Rrightarrow$ *body*$_2$ $\Longrightarrow$
(Γ, τ) $\vdash_C$ *Iterator iter its None body*$_1$ $\Rrightarrow$ *Iterator iter its* (*Some* σ) *body*$_2$

|*ExprListNilN*:
$\Gamma \vdash_L [] \Rrightarrow []$
|*ExprListConsN*:
$\Gamma \vdash x \Rrightarrow y \Longrightarrow$
$\Gamma \vdash_L xs \Rrightarrow ys \Longrightarrow$
$\Gamma \vdash_L x \# xs \Rrightarrow y \# ys$

7.2 Elimination Rules

inductive-cases *LiteralNE* [*elim*]: $\Gamma \vdash$ *Literal a* $\Rrightarrow$ *b*
inductive-cases *LetNE* [*elim*]: $\Gamma \vdash$ *Let v t init body* $\Rrightarrow$ *b*
inductive-cases *VarNE* [*elim*]: $\Gamma \vdash$ *Var v* $\Rrightarrow$ *b*
inductive-cases *IfNE* [*elim*]: $\Gamma \vdash$ *If a b c* $\Rrightarrow$ *d*

inductive-cases *MetaOperationCallNE* [*elim*]: $\Gamma \vdash$ *MetaOperationCall* τ *op* $\Rrightarrow$ *b*

inductive-cases *StaticOperationCallNE* [*elim*]: $\Gamma \vdash$ *StaticOperationCall* τ *op as* $\Rrightarrow$ *b*
inductive-cases *DotCallNE* [*elim*]: $\Gamma \vdash$ *Call src DotCall call* $\Rrightarrow$ *b*
inductive-cases *SafeDotCallNE* [*elim*]: $\Gamma \vdash$ *Call src SafeDotCall call* $\Rrightarrow$ *b*
inductive-cases *ArrowCallNE* [*elim*]: $\Gamma \vdash$ *Call src ArrowCall call* $\Rrightarrow$ *b*
inductive-cases *SafeArrowCallNE* [*elim*]: $\Gamma \vdash$ *Call src SafeArrowCall call* $\Rrightarrow$ *b*

inductive-cases *CallNE* [*elim*]: (Γ, τ) $\vdash_C$ *call* $\Rrightarrow$ *b*
inductive-cases *OperationCallNE* [*elim*]: (Γ, τ) $\vdash_C$ *Operation op as* $\Rrightarrow$ *call*
inductive-cases *IterateCallNE* [*elim*]: (Γ, τ) $\vdash_C$ *Iterate its its-ty res res-t res-init body* $\Rrightarrow$ *call*
inductive-cases *IteratorCallNE* [*elim*]: (Γ, τ) $\vdash_C$ *Iterator iter its its-ty body* $\Rrightarrow$ *call*

inductive-cases *ExprListNE* [*elim*]: $\Gamma \vdash_L$ *xs* $\Rrightarrow$ *ys*

7.3 Simplification Rules

inductive-simps *normalize-alt-simps*:
$\Gamma \vdash \mathit{Literal}\ a \Rrightarrow b$
$\Gamma \vdash \mathit{Let}\ v\ t\ \mathit{init}\ \mathit{body} \Rrightarrow b$
$\Gamma \vdash \mathit{Var}\ v \Rrightarrow b$
$\Gamma \vdash \mathit{If}\ a\ b\ c \Rrightarrow d$

$\Gamma \vdash \mathit{MetaOperationCall}\ \tau\ \mathit{op} \Rrightarrow b$
$\Gamma \vdash \mathit{StaticOperationCall}\ \tau\ \mathit{op}\ \mathit{as} \Rrightarrow b$
$\Gamma \vdash \mathit{Call}\ \mathit{src}\ \mathit{DotCall}\ \mathit{call} \Rrightarrow b$
$\Gamma \vdash \mathit{Call}\ \mathit{src}\ \mathit{SafeDotCall}\ \mathit{call} \Rrightarrow b$
$\Gamma \vdash \mathit{Call}\ \mathit{src}\ \mathit{ArrowCall}\ \mathit{call} \Rrightarrow b$
$\Gamma \vdash \mathit{Call}\ \mathit{src}\ \mathit{SafeArrowCall}\ \mathit{call} \Rrightarrow b$

$(\Gamma, \tau) \vdash_C \mathit{call} \Rrightarrow b$
$(\Gamma, \tau) \vdash_C \mathit{Operation}\ \mathit{op}\ \mathit{as} \Rrightarrow \mathit{call}$
$(\Gamma, \tau) \vdash_C \mathit{Iterate}\ \mathit{its}\ \mathit{its\text{-}ty}\ \mathit{res}\ \mathit{res\text{-}t}\ \mathit{res\text{-}init}\ \mathit{body} \Rrightarrow \mathit{call}$
$(\Gamma, \tau) \vdash_C \mathit{Iterator}\ \mathit{iter}\ \mathit{its}\ \mathit{its\text{-}ty}\ \mathit{body} \Rrightarrow \mathit{call}$

$\Gamma \vdash_L [] \Rrightarrow ys$
$\Gamma \vdash_L x \# xs \Rrightarrow ys$

7.4 Determinism

lemma *any-has-not-element-type*:
$\mathit{element\text{-}type}\ \tau\ \sigma \Longrightarrow \tau \leq \mathit{OclAny}[?] \vee \tau \leq \mathit{Tuple}\ \mathit{fmempty} \Longrightarrow \mathit{False}$
by (*erule element-type.cases*; *auto*)

lemma *any-has-not-element-type'*:
$\mathit{element\text{-}type}\ \tau\ \sigma \Longrightarrow \mathit{OclVoid}[?] \leq \tau \Longrightarrow \mathit{False}$
by (*erule element-type.cases*; *auto*)

lemma
normalize-det:
$\Gamma \vdash \mathit{expr} \Rrightarrow \mathit{expr}_1 \Longrightarrow$
$\Gamma \vdash \mathit{expr} \Rrightarrow \mathit{expr}_2 \Longrightarrow \mathit{expr}_1 = \mathit{expr}_2$ **and**
normalize-call-det:
$\Gamma\text{-}\tau \vdash_C \mathit{call} \Rrightarrow \mathit{call}_1 \Longrightarrow$
$\Gamma\text{-}\tau \vdash_C \mathit{call} \Rrightarrow \mathit{call}_2 \Longrightarrow \mathit{call}_1 = \mathit{call}_2$ **and**
normalize-expr-list-det:
$\Gamma \vdash_L xs \Rrightarrow ys \Longrightarrow$
$\Gamma \vdash_L xs \Rrightarrow zs \Longrightarrow ys = zs$
for Γ :: ('*a* :: *ocl-object-model*) *type env*
and $\Gamma\text{-}\tau$ :: ('*a* :: *ocl-object-model*) *type env* $\times$ '*a type*
proof (*induct arbitrary*: expr_2 **and** call_2 **and** *zs*
rule: *normalize-normalize-call-normalize-expr-list.inducts*)
case (*LiteralN* Γ *a*) **thus** *?case* **by** *auto*
next

```
  case (ExplicitlyTypedLetN Γ init₁ init₂ v τ body₁ body₂) thus ?case
    by blast
next
  case (ImplicitlyTypedLetN Γ init₁ init₂ τ v body₁ body₂) thus ?case
    by (metis (mono-tags, lifting) LetNE option.distinct(1) typing-det)
next
  case (VarN Γ v) thus ?case by auto
next
  case (IfN Γ a₁ a₂ b₁ b₂ c₁ c₂) thus ?case
    apply (insert IfN.prems)
    apply (erule IfNE)
    by (simp add: IfN.hyps)
next
  case (MetaOperationCallN Γ τ op) thus ?case by auto
next
  case (StaticOperationCallN Γ params₁ params₂ τ op) thus ?case by blast
next
  case (OclAnyDotCallN Γ src₁ src₂ τ call₁ call₂) show ?case
    apply (insert OclAnyDotCallN.prems)
    apply (erule DotCallNE)
    using OclAnyDotCallN.hyps typing-det apply metis
    using OclAnyDotCallN.hyps any-has-not-element-type typing-det by metis
next
  case (OclAnySafeDotCallN Γ src₁ src₂ τ call₁ call₂) show ?case
    apply (insert OclAnySafeDotCallN.prems)
    apply (erule SafeDotCallNE)
    using OclAnySafeDotCallN.hyps typing-det comp-apply
    apply (metis (no-types, lifting) list.simps(8) list.simps(9))
    using OclAnySafeDotCallN.hyps typing-det any-has-not-element-type'
    by metis
next
  case (OclAnyArrowCallN Γ src₁ src₂ τ src₃ σ call₁ call₂) show ?case
    apply (insert OclAnyArrowCallN.prems)
    apply (erule ArrowCallNE)
    using OclAnyArrowCallN.hyps typing-det comp-apply apply metis
    using OclAnyArrowCallN.hyps typing-det any-has-not-element-type
    by metis
next
  case (CollectionArrowCallN Γ src₁ src₂ τ uu call₁ call₂) show ?case
    apply (insert CollectionArrowCallN.prems)
    apply (erule ArrowCallNE)
    using CollectionArrowCallN.hyps typing-det any-has-not-element-type
    apply metis
    using CollectionArrowCallN.hyps typing-det by metis
next
  case (CollectionSafeArrowCallN Γ src₁ src₂ τ σ src₃ ϱ call₁ call₂) show ?case
    apply (insert CollectionSafeArrowCallN.prems)
    apply (erule SafeArrowCallNE)
    using CollectionSafeArrowCallN.hyps typing-det element-type-det by metis
```

next
case (*CollectionDotCallN* Γ src_1 src_2 τ σ $call_1$ $call_2$ *it*) **show** *?case*
apply (*insert CollectionDotCallN.prems*)
apply (*erule DotCallNE*)
using *CollectionDotCallN.hyps typing-det any-has-not-element-type*
apply *metis*
using *CollectionDotCallN.hyps typing-det element-type-det* **by** *metis*
next
case (*CollectionSafeDotCallN* Γ src_1 src_2 τ σ src_3 $call_1$ $call_2$ *it*) **show** *?case*
apply (*insert CollectionSafeDotCallN.prems*)
apply (*erule SafeDotCallNE*)
using *CollectionSafeDotCallN.hyps typing-det any-has-not-element-type'*
apply *metis*
using *CollectionSafeDotCallN.hyps typing-det element-type-det* **by** *metis*
next
case (*TypeOperationN* Γ τ *op ty*) **thus** *?case* **by** *auto*
next
case (*AttributeN* Γ τ *attr*) **thus** *?case* **by** *auto*
next
case (*AssociationEndN* Γ τ *role from*) **thus** *?case* **by** *auto*
next
case (*AssociationClassN* Γ τ $\mathcal{A}$ *from*) **thus** *?case* **by** *auto*
next
case (*AssociationClassEndN* Γ τ *role*) **thus** *?case* **by** *auto*
next
case (*OperationN* Γ $params_1$ $params_2$ τ *op*) **thus** *?case* **by** *blast*
next
case (*TupleElementN* Γ τ *elem*) **thus** *?case* **by** *auto*
next
case (*ExplicitlyTypedIterateN*
Γ $res\text{-}init_1$ $res\text{-}init_2$ σ *its res* $res\text{-}t_1$ $body_1$ $res\text{-}t_2$ $body_2$ τ)
show *?case*
apply (*insert ExplicitlyTypedIterateN.prems*)
apply (*erule IterateCallNE*)
using *ExplicitlyTypedIterateN.hyps element-type-det* **by** *blast+*
next
case (*ImplicitlyTypedIterateN*
τ σ Γ $res\text{-}init_1$ $res\text{-}init_2$ *its res* $res\text{-}t_1$ $body_1$ $res\text{-}t_2$ $body_2$)
show *?case*
apply (*insert ImplicitlyTypedIterateN.prems*)
apply (*erule IterateCallNE*)
using *ImplicitlyTypedIterateN.hyps element-type-det* **by** *blast+*
next
case (*ExplicitlyTypedIteratorN* Γ σ *its* $body_1$ $body_2$ τ *iter*)
show *?case*
apply (*insert ExplicitlyTypedIteratorN.prems*)
apply (*erule IteratorCallNE*)
using *ExplicitlyTypedIteratorN.hyps element-type-det* **by** *blast+*
next

case (*ImplicitlyTypedIteratorN* τ σ Γ *its* $body_1$ $body_2$ *iter*)
show *?case*
apply (*insert ImplicitlyTypedIteratorN.prems*)
apply (*erule IteratorCallNE*)
using *ImplicitlyTypedIteratorN.hyps element-type-det* **by** *blast+*
next
case (*ExprListNilN* Γ) **thus** *?case*
using *normalize-expr-list.cases* **by** *auto*
next
case (*ExprListConsN* Γ *x y xs ys*) **thus** *?case* **by** *blast*
qed

7.5 Normalized Expressions Typing

Here is the final typing rules.

inductive *nf-typing* (‹(*1-*/ $\vdash$/ (- :/ -))› [*51*,*51*,*51*] *50*) **where**
$\Gamma \vdash expr \Rrightarrow expr_N \Longrightarrow$
$\Gamma \vdash_E expr_N : \tau \Longrightarrow$
$\Gamma \vdash expr : \tau$

lemma *nf-typing-det*:
$\Gamma \vdash expr : \tau \Longrightarrow$
$\Gamma \vdash expr : \sigma \Longrightarrow \tau = \sigma$
by (*metis nf-typing.cases normalize-det typing-det*)

7.6 Code Setup

code-pred *normalize* **.**

code-pred *nf-typing* **.**

definition *check-type* Γ *expr* $\tau \equiv$
Predicate.eval (*nf-typing-i-i-i* Γ *expr* τ) ()

definition *synthesize-type* Γ *expr* $\equiv$
Predicate.singleton (λ-. *OclInvalid*)
(*Predicate.map errorable* (*nf-typing-i-i-o* Γ *expr*))

It is the only usage of the *OclInvalid* type. This type is not required to define typing rules. It is only required to make the typing function total.

end

Chapter 8

Examples

theory *OCL-Examples*
  imports *OCL-Normalization*
begin

8.1 Classes

datatype *classes1* =
  Object | *Person* | *Employee* | *Customer* | *Project* | *Task* | *Sprint*

inductive *subclass1* **where**
   $c \neq Object \Longrightarrow$
   subclass1 c Object
| *subclass1 Employee Person*
| *subclass1 Customer Person*

instantiation *classes1* :: *semilattice-sup*
begin

definition $(<) \equiv$ *subclass1*
definition $(\leq) \equiv$ *subclass1*$^{==}$

fun *sup-classes1* **where**
   Object $\sqcup$ - = *Object*
| *Person* $\sqcup$ *c* = (*if c* = *Person* $\vee$ *c* = *Employee* $\vee$ *c* = *Customer*
   then Person else Object)
| *Employee* $\sqcup$ *c* = (*if c* = *Employee then Employee else*
   if c = *Person* $\vee$ *c* = *Customer then Person else Object*)
| *Customer* $\sqcup$ *c* = (*if c* = *Customer then Customer else*
   if c = *Person* $\vee$ *c* = *Employee then Person else Object*)
| *Project* $\sqcup$ *c* = (*if c* = *Project then Project else Object*)
| *Task* $\sqcup$ *c* = (*if c* = *Task then Task else Object*)
| *Sprint* $\sqcup$ *c* = (*if c* = *Sprint then Sprint else Object*)

lemma *less-le-not-le-classes1*:

$c < d \longleftrightarrow c \leq d \wedge \neg\, d \leq c$
for $c\ d$ :: *classes1*
unfolding *less-classes1-def less-eq-classes1-def*
using *subclass1.simps* **by** *auto*

lemma *order-refl-classes1*:
$c \leq c$
for c :: *classes1*
unfolding *less-eq-classes1-def* **by** *simp*

lemma *order-trans-classes1*:
$c \leq d \Longrightarrow d \leq e \Longrightarrow c \leq e$
for $c\ d\ e$ :: *classes1*
unfolding *less-eq-classes1-def*
using *subclass1.simps* **by** *auto*

lemma *antisym-classes1*:
$c \leq d \Longrightarrow d \leq c \Longrightarrow c = d$
for $c\ d$ :: *classes1*
unfolding *less-eq-classes1-def*
using *subclass1.simps* **by** *auto*

lemma *sup-ge1-classes1*:
$c \leq c \sqcup d$
for $c\ d$ :: *classes1*
by (*induct c*; *auto simp add*: *less-eq-classes1-def less-classes1-def subclass1.simps*)

lemma *sup-ge2-classes1*:
$d \leq c \sqcup d$
for $c\ d$ :: *classes1*
by (*induct c*; *auto simp add*: *less-eq-classes1-def less-classes1-def subclass1.simps*)

lemma *sup-least-classes1*:
$c \leq e \Longrightarrow d \leq e \Longrightarrow c \sqcup d \leq e$
for $c\ d\ e$ :: *classes1*
by (*induct c*; *induct d*;
auto simp add: *less-eq-classes1-def less-classes1-def subclass1.simps*)

instance
apply *intro-classes*
apply (*simp add*: *less-le-not-le-classes1*)
apply (*simp add*: *order-refl-classes1*)
apply (*rule order-trans-classes1*; *auto*)
apply (*simp add*: *antisym-classes1*)
apply (*simp add*: *sup-ge1-classes1*)
apply (*simp add*: *sup-ge2-classes1*)
by (*simp add*: *sup-least-classes1*)

end

code-pred *subclass1* **.**

fun *subclass1-fun* **where**
  subclass1-fun Object $\mathcal{C}$ = *False*
| *subclass1-fun Person* $\mathcal{C}$ = ($\mathcal{C}$ = *Object*)
| *subclass1-fun Employee* $\mathcal{C}$ = ($\mathcal{C}$ = *Object* ∨ $\mathcal{C}$ = *Person*)
| *subclass1-fun Customer* $\mathcal{C}$ = ($\mathcal{C}$ = *Object* ∨ $\mathcal{C}$ = *Person*)
| *subclass1-fun Project* $\mathcal{C}$ = ($\mathcal{C}$ = *Object*)
| *subclass1-fun Task* $\mathcal{C}$ = ($\mathcal{C}$ = *Object*)
| *subclass1-fun Sprint* $\mathcal{C}$ = ($\mathcal{C}$ = *Object*)

lemma *less-classes1-code* [*code*]:
  (<) = *subclass1-fun*
proof (*intro ext iffI*)
 fix $\mathcal{C}$ $\mathcal{D}$ :: *classes1*
 show $\mathcal{C}$ < $\mathcal{D}$ ⟹ *subclass1-fun* $\mathcal{C}$ $\mathcal{D}$
  unfolding *less-classes1-def*
  apply (*erule subclass1.cases*, *auto*)
  using *subclass1-fun.elims*(*3*) **by** *blast*
 show *subclass1-fun* $\mathcal{C}$ $\mathcal{D}$ ⟹ $\mathcal{C}$ < $\mathcal{D}$
  by (*erule subclass1-fun.elims*, *auto simp add*: *less-classes1-def subclass1.intros*)
qed

lemma *less-eq-classes1-code* [*code*]:
  (≤) = (λ*x y*. *subclass1-fun x y* ∨ *x* = *y*)
 unfolding *dual-order.order-iff-strict less-classes1-code*
 by *auto*

8.2 Object Model

abbreviation Γ_0 ≡ *fmempty* :: *classes1 type env*
declare [[*coercion ObjectType* :: *classes1* ⇒ *classes1 basic-type*]]
declare [[*coercion phantom* :: *String.literal* ⇒ *classes1 enum*]]

instantiation *classes1* :: *ocl-object-model*
begin

definition *classes-classes1* ≡
 {|*Object*, *Person*, *Employee*, *Customer*, *Project*, *Task*, *Sprint*|}

definition *attributes-classes1* ≡ *fmap-of-list* [
 (*Person*, *fmap-of-list* [
  (*STR* ''*name*'', *String*[*1*] :: *classes1 type*)]),
 (*Employee*, *fmap-of-list* [
  (*STR* ''*name*'', *String*[*1*]),
  (*STR* ''*position*'', *String*[*1*])]),
 (*Customer*, *fmap-of-list* [
  (*STR* ''*vip*'', *Boolean*[*1*])]),

```
  (Project, fmap-of-list [
    (STR ''name'', String[1]),
    (STR ''cost'', Real[?])]),
  (Task, fmap-of-list [
    (STR ''description'', String[1])])]

abbreviation assocs ≡ [
  STR ''ProjectManager'' ↦f [
    STR ''projects'' ↦f (Project, 0::nat, ∞::enat, False, True),
    STR ''manager'' ↦f (Employee, 1, 1, False, False)],
  STR ''ProjectMember'' ↦f [
    STR ''member-of'' ↦f (Project, 0, ∞, False, False),
    STR ''members'' ↦f (Employee, 1, 20, True, True)],
  STR ''ManagerEmployee'' ↦f [
    STR ''line-manager'' ↦f (Employee, 0, 1, False, False),
    STR ''project-manager'' ↦f (Employee, 0, ∞, False, False),
    STR ''employees'' ↦f (Employee, 3, 7, False, False)],
  STR ''ProjectCustomer'' ↦f [
    STR ''projects'' ↦f (Project, 0, ∞, False, True),
    STR ''customer'' ↦f (Customer, 1, 1, False, False)],
  STR ''ProjectTask'' ↦f [
    STR ''project'' ↦f (Project, 1, 1, False, False),
    STR ''tasks'' ↦f (Task, 0, ∞, True, True)],
  STR ''SprintTaskAssignee'' ↦f [
    STR ''sprint'' ↦f (Sprint, 0, 10, False, True),
    STR ''tasks'' ↦f (Task, 0, 5, False, True),
    STR ''assignee'' ↦f (Employee, 0, 1, False, False)]]

definition associations-classes1 ≡ assocs

definition association-classes-classes1 ≡ fmempty :: classes1 ⇀f assoc


context Project
def: membersCount() : Integer[1] = members->size()
def: membersByName(mn : String[1]) : Set(Employee[1]) =
       members->select(member | member.name = mn)
static def: allProjects() : Set(Project[1]) =
              Project[1].allInstances()


definition operations-classes1 ≡ [
  (STR ''membersCount'', Project[1], [], Integer[1], False,
   Some (OperationCall
    (AssociationEndCall (Var STR ''self'') DotCall None STR ''members'')
    ArrowCall CollectionSizeOp [])),
  (STR ''membersByName'', Project[1], [(STR ''mn'', String[1], In)],
   Set Employee[1], False,
   Some (SelectIteratorCall
    (AssociationEndCall (Var STR ''self'') DotCall None STR ''members'')
```

```
    ArrowCall [STR ''member''] None
    (OperationCall
      (AttributeCall (Var STR ''member'') DotCall STR ''name'')
      DotCall EqualOp [Var STR ''mn'']))),
  (STR ''allProjects'', Project[1], [], Set Project[1], True,
   Some (MetaOperationCall Project[1] AllInstancesOp))
  ] :: (classes1 type, classes1 expr) oper-spec list

definition literals-classes1 ≡ fmap-of-list [
  (STR ''E1'' :: classes1 enum, {|STR ''A'', STR ''B''|}),
  (STR ''E2'', {|STR ''C'', STR ''D'', STR ''E''|})]


lemma assoc-end-min-less-eq-max:
    assoc |∈| fmdom assocs ⟹
    fmlookup assocs assoc = Some ends ⟹
    role |∈| fmdom ends ⟹
    fmlookup ends role = Some end ⟹
    assoc-end-min end ≤ assoc-end-max end
  unfolding assoc-end-min-def assoc-end-max-def
  using zero-enat-def one-enat-def numeral-eq-enat by auto

lemma association-ends-unique:
  assumes association-ends' classes assocs C from role end1
      and association-ends' classes assocs C from role end2
    shows end1 = end2
proof −
  have ¬ association-ends-not-unique' classes assocs by eval
  with assms show ?thesis
    using association-ends-not-unique'.simps by blast
qed

instance
  apply standard
  unfolding associations-classes1-def
  using assoc-end-min-less-eq-max apply blast
  using association-ends-unique by blast

end
```

8.3 Simplification Rules

```
lemma ex-alt-simps [simp]:
    ∃ a. a
    ∃ a. ¬ a
    (∃ a. (a ⟶ P) ∧ a) = P
    (∃ a. ¬ a ∧ (¬ a ⟶ P)) = P
  by auto
```

declare *numeral-eq-enat* [*simp*]

lemmas *basic-type-le-less* [*simp*] = *Orderings.order-class.le-less*
  for *x y* :: *'a basic-type*

declare *element-type-alt-simps* [*simp*]
declare *update-element-type.simps* [*simp*]
declare *to-unique-collection.simps* [*simp*]
declare *to-nonunique-collection.simps* [*simp*]
declare *to-ordered-collection.simps* [*simp*]

declare *assoc-end-class-def* [*simp*]
declare *assoc-end-min-def* [*simp*]
declare *assoc-end-max-def* [*simp*]
declare *assoc-end-ordered-def* [*simp*]
declare *assoc-end-unique-def* [*simp*]

declare *oper-name-def* [*simp*]
declare *oper-context-def* [*simp*]
declare *oper-params-def* [*simp*]
declare *oper-result-def* [*simp*]
declare *oper-static-def* [*simp*]
declare *oper-body-def* [*simp*]

declare *oper-in-params-def* [*simp*]
declare *oper-out-params-def* [*simp*]

declare *assoc-end-type-def* [*simp*]
declare *oper-type-def* [*simp*]

declare *op-type-alt-simps* [*simp*]
declare *typing-alt-simps* [*simp*]
declare *normalize-alt-simps* [*simp*]
declare *nf-typing.simps* [*simp*]

declare *subclass1.intros* [*intro*]
declare *less-classes1-def* [*simp*]

declare *literals-classes1-def* [*simp*]

lemma *attribute-Employee-name* [*simp*]:
  attribute Employee STR ''name'' $\mathcal{D}$ τ =
  ($\mathcal{D}$ = Employee $\wedge$ τ = String[1])
proof −
  have *attribute Employee STR ''name'' Employee String[1]*
    by *eval*
  thus *?thesis*
    using *attribute-det* **by** *blast*
qed

lemma *association-end-Project-members* [*simp*]:
association-end Project None STR ''members'' $\mathcal{D}$ τ =
($\mathcal{D}$ = Project $\wedge$ τ = (Employee, 1, 20, True, True))
proof −
have *association-end Project None STR ''members''*
Project (Employee, 1, 20, True, True)
by *eval*
thus *?thesis*
using *association-end-det* **by** *blast*
qed

lemma *association-end-Employee-projects-simp* [*simp*]:
association-end Employee None STR ''projects'' $\mathcal{D}$ τ =
($\mathcal{D}$ = Employee $\wedge$ τ = (Project, 0, ∞, False, True))
proof −
have *association-end Employee None STR ''projects''*
Employee (Project, 0, ∞, False, True)
by *eval*
thus *?thesis*
using *association-end-det* **by** *blast*
qed

lemma *static-operation-Project-allProjects* [*simp*]:
static-operation $\langle$Project$\rangle_{\mathcal{T}}$[1] STR ''allProjects'' [] oper =
(oper = (STR ''allProjects'', $\langle$Project$\rangle_{\mathcal{T}}$[1], [], Set $\langle$Project$\rangle_{\mathcal{T}}$[1], True,
Some (MetaOperationCall $\langle$Project$\rangle_{\mathcal{T}}$[1] AllInstancesOp)))
proof −
have *static-operation $\langle$Project$\rangle_{\mathcal{T}}$[1] STR ''allProjects'' []*
(STR ''allProjects'', $\langle$Project$\rangle_{\mathcal{T}}$[1], [], Set $\langle$Project$\rangle_{\mathcal{T}}$[1], True,
Some (MetaOperationCall $\langle$Project$\rangle_{\mathcal{T}}$[1] AllInstancesOp))
by *eval*
thus *?thesis*
using *static-operation-det* **by** *blast*
qed

8.4 Basic Types

8.4.1 Positive Cases

lemma *UnlimitedNatural* < (*Real* :: *classes1 basic-type*) **by** *simp*
lemma $\langle$*Employee*$\rangle_{\mathcal{T}}$ < $\langle$*Person*$\rangle_{\mathcal{T}}$ **by** *auto*
lemma $\langle$*Person*$\rangle_{\mathcal{T}}$ $\leq$ *OclAny* **by** *simp*

8.4.2 Negative Cases

lemma $\neg$ *String* $\leq$ (*Boolean* :: *classes1 basic-type*) **by** *simp*

8.5 Types

8.5.1 Positive Cases

lemma *Integer*[*?*] < (*OclSuper* :: *classes1 type*) **by** *simp*
lemma *Collection Real*[*?*] < (*OclSuper* :: *classes1 type*) **by** *simp*
lemma *Set* (*Collection Boolean*[*1*]) < (*OclSuper* :: *classes1 type*) **by** *simp*
lemma *Set* (*Bag Boolean*[*1*]) < *Set* (*Collection Boolean*[*?*] :: *classes1 type*)
by *simp*
lemma *Tuple* (*fmap-of-list* [(*STR* ''*a*'', *Boolean*[*1*]), (*STR* ''*b*'', *Integer*[*1*])]) <
Tuple (*fmap-of-list* [(*STR* ''*a*'', *Boolean*[*?*] :: *classes1 type*)]) **by** *eval*

lemma *Integer*[*1*] ⊔ (*Real*[*?*] :: *classes1 type*) = *Real*[*?*] **by** *simp*
lemma *Set Integer*[*1*] ⊔ *Set* (*Real*[*1*] :: *classes1 type*) = *Set Real*[*1*] **by** *simp*
lemma *Set Integer*[*1*] ⊔ *Bag* (*Boolean*[*?*] :: *classes1 type*) = *Collection OclAny*[*?*]

by *simp*
lemma *Set Integer*[*1*] ⊔ (*Real*[*1*] :: *classes1 type*) = *OclSuper* **by** *simp*

8.5.2 Negative Cases

lemma ¬ *OrderedSet Boolean*[*1*] < *Set* (*Boolean*[*1*] :: *classes1 type*) **by** *simp*

8.6 Typing

8.6.1 Positive Cases

```
E1::A : E1[1]
```

lemma
Γ$_0$ ⊢ *EnumLiteral STR* ''*E1*'' *STR* ''*A*'' : (*Enum STR* ''*E1*'')[*1*]
by *simp*

```
true or false : Boolean[1]
```

lemma
Γ$_0$ ⊢ *OperationCall* (*BooleanLiteral True*) *DotCall OrOp*
[*BooleanLiteral False*] : *Boolean*[*1*]
by *simp*

```
null and true : Boolean[?]
```

lemma
Γ$_0$ ⊢ *OperationCall* (*NullLiteral*) *DotCall AndOp*
[*BooleanLiteral True*] : *Boolean*[*?*]
by *simp*

```
let x : Real[1] = 5 in x + 7 : Real[1]
```

lemma
Γ$_0$ ⊢ *Let* (*STR* ''*x*'') (*Some Real*[*1*]) (*IntegerLiteral 5*)
(*OperationCall* (*Var STR* ''*x*'') *DotCall PlusOp* [*IntegerLiteral 7*]) : *Real*[*1*]
by *simp*

```
null.oclIsUndefined() : Boolean[1]
```

lemma
$\Gamma_0 \vdash$ *OperationCall* (*NullLiteral*) *DotCall OclIsUndefinedOp* [] : *Boolean*[*1*]
by *simp*

```
Sequence{1..5, null}.oclIsUndefined() : Sequence(Boolean[1])
```

lemma
$\Gamma_0 \vdash$ *OperationCall* (*CollectionLiteral SequenceKind*
[*CollectionRange* (*IntegerLiteral 1*) (*IntegerLiteral 5*),
CollectionItem NullLiteral])
DotCall OclIsUndefinedOp [] : *Sequence Boolean*[*1*]
by *simp*

```
Sequence{1..5}->product(Set{'a', 'b'})
  : Set(Tuple(first: Integer[1], second: String[1]))
```

lemma
$\Gamma_0 \vdash$ *OperationCall* (*CollectionLiteral SequenceKind*
[*CollectionRange* (*IntegerLiteral 1*) (*IntegerLiteral 5*)])
ArrowCall ProductOp
[*CollectionLiteral SetKind*
[*CollectionItem* (*StringLiteral ''a''*),
CollectionItem (*StringLiteral ''b''*)]] :
Set (*Tuple* (*fmap-of-list* [
(*STR ''first''*, *Integer*[*1*]), (*STR ''second''*, *String*[*1*])]))
by *simp*

```
Sequence{1..5, null}?->iterate(x, acc : Real[1] = 0 | acc + x)
  : Real[1]
```

lemma
$\Gamma_0 \vdash$ *IterateCall* (*CollectionLiteral SequenceKind*
[*CollectionRange* (*IntegerLiteral 1*) (*IntegerLiteral 5*),
CollectionItem NullLiteral]) *SafeArrowCall*
[*STR ''x''*] *None*
(*STR ''acc''*) (*Some Real*[*1*]) (*IntegerLiteral 0*)
(*OperationCall* (*Var STR ''acc''*) *DotCall PlusOp* [*Var STR ''x''*]) : *Real*[*1*]
by *simp*

```
Sequence{1..5, null}?->max() : Integer[1]
```

lemma
$\Gamma_0 \vdash$ *OperationCall* (*CollectionLiteral SequenceKind*
[*CollectionRange* (*IntegerLiteral 1*) (*IntegerLiteral 5*),
CollectionItem NullLiteral])
SafeArrowCall CollectionMaxOp [] : *Integer*[*1*]
by *simp*

```
let x : Sequence(String[?]) = Sequence{'abc', 'zxc'} in
x->any(it | it = 'test') : String[?]
```

lemma

$\Gamma_0 \vdash$ *Let* (*STR* ''*x*'') (*Some* (*Sequence String*[*?*]))
(*CollectionLiteral SequenceKind*
[*CollectionItem* (*StringLiteral* ''*abc*''),
CollectionItem (*StringLiteral* ''*zxc*'')])
(*AnyIteratorCall* (*Var STR* ''*x*'') *ArrowCall*
[*STR* ''*it*''] *None*
(*OperationCall* (*Var STR* ''*it*'') *DotCall EqualOp*
[*StringLiteral* ''*test*'')) : *String*[*?*]
by *simp*

```
let x : Sequence(String[?]) = Sequence{'abc', 'zxc'} in
x?->closure(it | it) : OrderedSet(String[1])
```

lemma
$\Gamma_0 \vdash$ *Let STR* ''*x*'' (*Some* (*Sequence String*[*?*]))
(*CollectionLiteral SequenceKind*
[*CollectionItem* (*StringLiteral* ''*abc*''),
CollectionItem (*StringLiteral* ''*zxc*'')])
(*ClosureIteratorCall* (*Var STR* ''*x*'') *SafeArrowCall*
[*STR* ''*it*''] *None*
(*Var STR* ''*it*'')) : *OrderedSet String*[*1*]
by *simp*

```
context Employee:
name : String[1]
```

lemma
Γ_0(*STR* ''*self*'' $\mapsto_f$ *Employee*[*1*]) $\vdash$
AttributeCall (*Var STR* ''*self*'') *DotCall STR* ''*name*'' : *String*[*1*]
by *simp*

```
context Employee:
projects : Set(Project[1])
```

lemma
Γ_0(*STR* ''*self*'' $\mapsto_f$ *Employee*[*1*]) $\vdash$
AssociationEndCall (*Var STR* ''*self*'') *DotCall None*
STR ''*projects*'' : *Set Project*[*1*]
by *simp*

```
context Employee:
projects.members : Bag(Employee[1])
```

lemma
Γ_0(*STR* ''*self*'' $\mapsto_f$ *Employee*[*1*]) $\vdash$
AssociationEndCall (*AssociationEndCall* (*Var STR* ''*self*'')
DotCall None STR ''*projects*'')
DotCall None STR ''*members*'' : *Bag Employee*[*1*]
by *simp*

```
Project[?].allInstances() : Set(Project[?])
```

lemma
$\Gamma_0 \vdash$ *MetaOperationCall Project*[*?*] *AllInstancesOp* : *Set Project*[*?*]

by *simp*

```
Project[1]::allProjects() : Set(Project[1])
```

lemma
$\Gamma_0 \vdash$ *StaticOperationCall Project*[*1*] *STR ''allProjects''* [] : *Set Project*[*1*]
by *simp*

8.6.2 Negative Cases

```
true = null
```

lemma
$\nexists \tau.\ \Gamma_0 \vdash$ *OperationCall* (*BooleanLiteral True*) *DotCall EqualOp*
[*NullLiteral*] : τ
by *simp*

```
let x : Boolean[1] = 5 in x and true
```

lemma
$\nexists \tau.\ \Gamma_0 \vdash$ *Let STR ''x''* (*Some Boolean*[*1*]) (*IntegerLiteral 5*)
(*OperationCall* (*Var STR ''x''*) *DotCall AndOp* [*BooleanLiteral True*]) : τ
by *simp*

```
let x : Sequence(String[?]) = Sequence{'abc', 'zxc'} in
x->closure(it | 1)
```

lemma
$\nexists \tau.\ \Gamma_0 \vdash$ *Let STR ''x''* (*Some* (*Sequence String*[*?*]))
(*CollectionLiteral SequenceKind*
[*CollectionItem* (*StringLiteral ''abc''*),
CollectionItem (*StringLiteral ''zxc''*)])
(*ClosureIteratorCall* (*Var STR ''x''*) *ArrowCall* [*STR ''it''*] *None*
(*IntegerLiteral 1*)) : τ
by *simp*

```
Sequence{1..5, null}->max()
```

lemma
$\nexists \tau.\ \Gamma_0 \vdash$ *OperationCall* (*CollectionLiteral SequenceKind*
[*CollectionRange* (*IntegerLiteral 1*) (*IntegerLiteral 5*),
CollectionItem NullLiteral])
ArrowCall CollectionMaxOp [] : τ
proof −
have ¬ *operation-defined* (*Integer*[*?*] :: *classes1 type*) *STR ''max''* [*Integer*[*?*]]
by *eval*
thus *?thesis* **by** *simp*
qed

8.7 Code

8.7.1 Positive Cases

values {($\mathcal{D}$, τ). *attribute Employee STR ''name''* $\mathcal{D}$ τ}

values {($\mathcal{D}$, *end*). *association-end Employee None STR ''employees'' $\mathcal{D}$ end*}
values {($\mathcal{D}$, *end*). *association-end Employee* (*Some STR ''project-manager''*) *STR ''employees'' $\mathcal{D}$ end*}
values {*op*. *operation Project*[*1*] *STR ''membersCount''* [] *op*}
values {*op*. *operation Project*[*1*] *STR ''membersByName''* [*String*[*1*]] *op*}
value *has-literal STR ''E1'' STR ''A''*

```
context Employee:
projects.members : Bag(Employee[1])
```

values
{τ. Γ_0(*STR ''self'' $\mapsto_f$ Employee*[*1*]) $\vdash$
AssociationEndCall (*AssociationEndCall* (*Var STR ''self''*)
DotCall None STR ''projects'')
DotCall None STR ''members'' : τ}

8.7.2 Negative Cases

values {($\mathcal{D}$, τ). *attribute Employee STR ''name2'' $\mathcal{D}$ τ*}
value *has-literal STR ''E1'' STR ''C''*

```
Sequence{1..5, null}->max()
```

values
{τ. Γ_0 $\vdash$ *OperationCall* (*CollectionLiteral SequenceKind*
[*CollectionRange* (*IntegerLiteral 1*) (*IntegerLiteral 5*),
CollectionItem NullLiteral])
ArrowCall CollectionMaxOp [] : τ}

end

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