

Roth’s Theorem on Arithmetic Progressions

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Abstract

We formalise a proof of Roth’s Theorem on Arithmetic Progressions, a major result in additive combinatorics on the existence of 3-term arithmetic progressions in subsets of natural numbers. To this end, we follow a proof using graph regularity. We employ our recent formalisation of Szemerédi’s Regularity Lemma, a major result in extremal graph theory, which we use here to prove the Triangle Counting Lemma and the Triangle Removal Lemma. Our sources are Yufei Zhao’s MIT lecture notes “Graph Theory and Additive Combinatorics”¹ and W.T. Gowers’s Cambridge lecture notes “Topics in Combinatorics”.² We also refer to the University of Georgia notes by Stephanie Bell and Will Grodzicki “Using Szemerédi’s Regularity Lemma to Prove Roth’s Theorem”.³

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¹<https://yufeizhao.com/gtacbook/> and <https://yufeizhao.com/gtac/gtac.pdf>

²<https://www.dpmms.cam.ac.uk/~par31/notes/tic.pdf>

³<http://citeseerx.ist.psu.edu/viewdoc/summary?doi=10.1.1.432.327>

1 Roth's Theorem on Arithmetic Progressions

theory *Roth-Arithmetic-Progressions*

imports *Szemerédi-Regularity.Szemerédi*
Random-Graph-Subgraph-Threshold.Subgraph-Threshold
Ergodic-Theory.Asymptotic-Density
HOL-Library.Ramsey HOL-Library.Nat-Bijection

begin

1.1 Miscellaneous Preliminaries

lemma *sum-prod-le-prod-sum*:

fixes $a :: 'a \Rightarrow 'b :: \text{linordered-idom}$

assumes $\bigwedge i. i \in I \implies a\ i \geq 0 \wedge b\ i \geq 0$

shows $(\sum i \in I. \sum j \in I. a\ i * b\ j) \leq (\sum i \in I. a\ i) * (\sum i \in I. b\ i)$
<proof>

lemma *real-mult-gt-cube*: $A \geq (X :: \text{real}) \implies B \geq X \implies C \geq X \implies X \geq 0 \implies A * B * C \geq X^3$

<proof>

lemma *triple-sigma-rewrite-card*:

assumes *finite X finite Y finite Z*

shows $\text{card } \{(x,y,z) . x \in X \wedge (y,z) \in Y \times Z \wedge P\ x\ y\ z\} = (\sum x \in X . \text{card } \{(y,z) \in Y \times Z . P\ x\ y\ z\})$
<proof>

lemma *all-edges-between-mono1*:

$Y \subseteq Z \implies \text{all-edges-between } Y\ X\ G \subseteq \text{all-edges-between } Z\ X\ G$

<proof>

lemma *all-edges-between-mono2*:

$Y \subseteq Z \implies \text{all-edges-between } X\ Y\ G \subseteq \text{all-edges-between } X\ Z\ G$

<proof>

lemma *uwellformed-alt-fst*:

assumes *uwellformed G {x, y} \in uedges G*

shows $x \in \text{uverts } G$

<proof>

lemma *uwellformed-alt-snd*:

assumes *uwellformed G {x, y} \in uedges G*

shows $y \in \text{uverts } G$

<proof>

lemma *all-edges-between-subset-times*: $\text{all-edges-between } X\ Y\ G \subseteq (X \cap \bigcup (\text{uedges } G)) \times (Y \cap \bigcup (\text{uedges } G))$

<proof>

lemma *finite-all-edges-between'*:
assumes *finite* (*uverts* G) *uwellformed* G
shows *finite* (*all-edges-between* X Y G)
<proof>

lemma *all-edges-between-E-diff*:
all-edges-between X Y ($V, E - E'$) = *all-edges-between* X Y (V, E) - *all-edges-between* X Y (V, E')
<proof>

lemma *all-edges-between-E-Un*:
all-edges-between X Y ($V, E \cup E'$) = *all-edges-between* X Y (V, E) \cup *all-edges-between* X Y (V, E')
<proof>

lemma *all-edges-between-E-UN*:
all-edges-between X Y ($V, \bigcup_{i \in I}. E_i$) = ($\bigcup_{i \in I}. \text{all-edges-between } X \ Y \ (V, E_i)$)
<proof>

lemma *all-edges-preserved*: $\llbracket \text{all-edges-between } A \ B \ G' = \text{all-edges-between } A \ B \ G; X \subseteq A; Y \subseteq B \rrbracket$
 $\implies \text{all-edges-between } X \ Y \ G' = \text{all-edges-between } X \ Y \ G$
<proof>

lemma *subgraph-edge-wf*:
assumes *uwellformed* G *uverts* $H = \text{uverts } G$ *uedges* $H \subseteq \text{uedges } G$
shows *uwellformed* H
<proof>

1.2 Preliminaries on Neighbors in Graphs

definition *neighbor-in-graph*:: *uvert* \Rightarrow *uvert* \Rightarrow *ugraph* \Rightarrow *bool*
where *neighbor-in-graph* $x \ y \ G \equiv (x \in (\text{uverts } G) \wedge y \in (\text{uverts } G) \wedge \{x, y\} \in (\text{uedges } G))$

definition *neighbors* :: *uvert* \Rightarrow *ugraph* \Rightarrow *uvert set* **where**
neighbors $x \ G \equiv \{y \in \text{uverts } G . \text{neighbor-in-graph } x \ y \ G\}$

definition *neighbors-ss*:: *uvert set* \Rightarrow *ugraph* \Rightarrow *uvert set* **where**
neighbors-ss $x \ Y \ G \equiv \{y \in Y . \text{neighbor-in-graph } x \ y \ G\}$

lemma *all-edges-betw-sigma-neighbor*:
uwellformed $G \implies \text{all-edges-between } X \ Y \ G = (\text{SIGMA } x:X. \text{neighbors-ss } x \ Y \ G)$
<proof>

lemma *card-all-edges-betw-neighbor*:
assumes *finite* X *finite* Y *uwellformed* G
shows *card* (*all-edges-between* $X \ Y \ G$) = ($\sum_{x \in X}. \text{card} (\text{neighbors-ss } x \ Y \ G)$)

<proof>

1.3 Preliminaries on Triangles in Graphs

definition *triangle-in-graph*:: $uvert \Rightarrow uvert \Rightarrow uvert \Rightarrow ugraph \Rightarrow bool$
where *triangle-in-graph* $x\ y\ z\ G \equiv (\{x,y\} \in uedges\ G) \wedge (\{y,z\} \in uedges\ G) \wedge (\{x,z\} \in uedges\ G)$

definition *triangle-triples*
where *triangle-triples* $X\ Y\ Z\ G \equiv \{(x,y,z) \in X \times Y \times Z. \text{triangle-in-graph } x\ y\ z\ G\}$

lemma *triangle-commu1*:
assumes *triangle-in-graph* $x\ y\ z\ G$
shows *triangle-in-graph* $y\ x\ z\ G$
<proof>

lemma *triangle-vertices-distinct1*:
assumes *wf*: *uwellformed* G
assumes *tri*: *triangle-in-graph* $x\ y\ z\ G$
shows $x \neq y$
<proof>

lemma *triangle-vertices-distinct2*:
assumes *uwellformed* G *triangle-in-graph* $x\ y\ z\ G$
shows $y \neq z$
<proof>

lemma *triangle-in-graph-edge-point*:
assumes *uwellformed* G
shows *triangle-in-graph* $x\ y\ z\ G \iff \{y, z\} \in uedges\ G \wedge \text{neighbor-in-graph } x\ y\ z\ G \wedge \text{neighbor-in-graph } x\ z\ G$
<proof>

definition
unique-triangles G
 $\equiv \forall e \in uedges\ G. \exists! T. \exists x\ y\ z. T = \{x,y,z\} \wedge \text{triangle-in-graph } x\ y\ z\ G \wedge e \subseteq T$

definition *triangle-free-graph*:: $ugraph \Rightarrow bool$
where *triangle-free-graph* $G \equiv \neg(\exists x\ y\ z. \text{triangle-in-graph } x\ y\ z\ G)$

lemma *triangle-free-graph-empty*: $uedges\ G = \{\} \implies \text{triangle-free-graph } G$
<proof>

lemma *edge-vertices-not-equal*:
assumes *uwellformed* G $\{x,y\} \in uedges\ G$
shows $x \neq y$
<proof>

lemma *triangle-in-graph-verts*:

assumes *uwellformed* G *triangle-in-graph* $x\ y\ z\ G$

shows $x \in \text{uverts } G$ $y \in \text{uverts } G$ $z \in \text{uverts } G$

<proof>

definition *triangle-set* :: *ugraph* \Rightarrow *uvert set set*

where *triangle-set* $G \equiv \{ \{x,y,z\} \mid x\ y\ z.\ \text{triangle-in-graph } x\ y\ z\ G \}$

fun *mk-triangle-set* :: (*uvert* \times *uvert* \times *uvert*) \Rightarrow *uvert set*

where *mk-triangle-set* $(x,y,z) = \{x,y,z\}$

lemma *finite-triangle-set*:

assumes *fin*: *finite* (*uverts* G) **and** *wf*: *uwellformed* G

shows *finite* (*triangle-set* G)

<proof>

lemma *card-triangle-3*:

assumes $t \in \text{triangle-set } G$ *uwellformed* G

shows *card* $t = 3$

<proof>

lemma *triangle-set-power-set-ss*: *uwellformed* $G \Longrightarrow \text{triangle-set } G \subseteq \text{Pow } (\text{uverts } G)$

<proof>

lemma *triangle-in-graph-ss*:

assumes *uedges* $G_{\text{new}} \subseteq \text{uedges } G$

assumes *triangle-in-graph* $x\ y\ z\ G_{\text{new}}$

shows *triangle-in-graph* $x\ y\ z\ G$

<proof>

lemma *triangle-set-graph-edge-ss*:

assumes *uedges* $G_{\text{new}} \subseteq \text{uedges } G$

assumes *uverts* $G_{\text{new}} = \text{uverts } G$

shows *triangle-set* $G_{\text{new}} \subseteq \text{triangle-set } G$

<proof>

lemma *triangle-set-graph-edge-ss-bound*:

fixes $G :: \text{ugraph}$ **and** $G_{\text{new}} :: \text{ugraph}$

assumes *uwellformed* G *finite* (*uverts* G) *uedges* $G_{\text{new}} \subseteq \text{uedges } G$ *uverts* $G_{\text{new}} = \text{uverts } G$

shows *card* (*triangle-set* G) \geq *card* (*triangle-set* G_{new})

<proof>

1.4 The Triangle Counting Lemma and the Triangle Removal Lemma

We begin with some more auxiliary material to be used in the main lemmas.

lemma *regular-pair-neighbor-bound*:

fixes $\varepsilon::\text{real}$
assumes $\text{fin}G$: *finite* (*uverts* G)
assumes xss : $X \subseteq \text{uverts } G$ **and** yss : $Y \subseteq \text{uverts } G$ **and** $\text{card } X > 0$
and wf : *uwellformed* G
and eg0 : $\varepsilon > 0$ **and** ε -*regular-pair* $X Y G$ **and** ed : *edge-density* $X Y G \geq 2*\varepsilon$
defines $X' \equiv \{x \in X. \text{card } (\text{neighbors-ss } x Y G) < (\text{edge-density } X Y G - \varepsilon) * \text{card } (Y)\}$
shows $\text{card } X' < \varepsilon * \text{card } X$
(is $\text{card } (?X') < \varepsilon * -$)
<proof>

lemma *neighbor-set-meets-e-reg-cond*:

fixes $\varepsilon::\text{real}$
assumes *edge-density* $X Y G \geq 2*\varepsilon$
and $\text{card } (\text{neighbors-ss } x Y G) \geq (\text{edge-density } X Y G - \varepsilon) * \text{card } Y$
shows $\text{card } (\text{neighbors-ss } x Y G) \geq \varepsilon * \text{card } (Y)$
<proof>

lemma *all-edges-btwn-neighbor-sets-lower-bound*:

fixes $\varepsilon::\text{real}$
assumes rp2 : ε -*regular-pair* $Y Z G$
and ed1 : *edge-density* $X Y G \geq 2*\varepsilon$ **and** ed2 : *edge-density* $X Z G \geq 2*\varepsilon$
and cond1 : $\text{card } (\text{neighbors-ss } x Y G) \geq (\text{edge-density } X Y G - \varepsilon) * \text{card } Y$
and cond2 : $\text{card } (\text{neighbors-ss } x Z G) \geq (\text{edge-density } X Z G - \varepsilon) * \text{card } Z$
shows $\text{card } (\text{all-edges-between } (\text{neighbors-ss } x Y G) (\text{neighbors-ss } x Z G) G)$
 $\geq (\text{edge-density } Y Z G - \varepsilon) * \text{card } (\text{neighbors-ss } x Y G) * \text{card } (\text{neighbors-ss } x Z G)$
(is $\text{card } (\text{all-edges-between } ?Y' ?Z' G) \geq (\text{edge-density } Y Z G - \varepsilon) * \text{card } ?Y'$
 $* \text{card } ?Z')$
<proof>

We are now ready to show the Triangle Counting Lemma:

theorem *triangle-counting-lemma*:

fixes $\varepsilon::\text{real}$
assumes xss : $X \subseteq \text{uverts } G$ **and** yss : $Y \subseteq \text{uverts } G$ **and** zss : $Z \subseteq \text{uverts } G$ **and**
 en0 : $\varepsilon > 0$
and $\text{fin}G$: *finite* (*uverts* G) **and** wf : *uwellformed* G
and rp1 : ε -*regular-pair* $X Y G$ **and** rp2 : ε -*regular-pair* $Y Z G$ **and** rp3 :
 ε -*regular-pair* $X Z G$
and ed1 : *edge-density* $X Y G \geq 2*\varepsilon$ **and** ed2 : *edge-density* $X Z G \geq 2*\varepsilon$ **and**
 ed3 : *edge-density* $Y Z G \geq 2*\varepsilon$
shows $\text{card } (\text{triangle-triples } X Y Z G)$
 $\geq (1-2*\varepsilon) * (\text{edge-density } X Y G - \varepsilon) * (\text{edge-density } X Z G - \varepsilon) * (\text{edge-density } Y Z G - \varepsilon)$

$\text{card } X * \text{card } Y * \text{card } Z$
 <proof>

definition *regular-graph* :: *real* \Rightarrow *uvert set set* \Rightarrow *ugraph* \Rightarrow *bool*
 (<--regular'-graph> [999]1000)
where *ε -regular-graph* $P G \equiv \forall R S. R \in P \longrightarrow S \in P \longrightarrow \varepsilon$ -regular-pair $R S G$
for $\varepsilon :: \text{real}$

A minimum density, but empty edge sets are excluded.

definition *edge-dense* :: *nat set* \Rightarrow *nat set* \Rightarrow *ugraph* \Rightarrow *real* \Rightarrow *bool*
where *edge-dense* $X Y G \varepsilon \equiv \text{all-edges-between } X Y G = \{\}$ \vee *edge-density* $X Y G \geq \varepsilon$

definition *dense-graph* :: *uvert set set* \Rightarrow *ugraph* \Rightarrow *real* \Rightarrow *bool*
where *dense-graph* $P G \varepsilon \equiv \forall R S. R \in P \longrightarrow S \in P \longrightarrow \text{edge-dense } R S G \varepsilon$ **for**
 $\varepsilon :: \text{real}$

definition *decent* :: *nat set* \Rightarrow *nat set* \Rightarrow *ugraph* \Rightarrow *real* \Rightarrow *bool*
where *decent* $X Y G \eta \equiv \text{all-edges-between } X Y G = \{\}$ \vee ($\text{card } X \geq \eta \wedge \text{card } Y \geq \eta$) **for** $\eta :: \text{real}$

definition *decent-graph* :: *uvert set set* \Rightarrow *ugraph* \Rightarrow *real* \Rightarrow *bool*
where *decent-graph* $P G \eta \equiv \forall R S. R \in P \longrightarrow S \in P \longrightarrow \text{decent } R S G \eta$

The proof of the triangle counting lemma requires ordered triples. For each unordered triple there are six permutations, hence the factor of 1/6 here.

lemma *card-convert-triangle-rep*:
fixes $G :: \text{ugraph}$
assumes $X \subseteq \text{uverts } G$ **and** $Y \subseteq \text{uverts } G$ **and** $Z \subseteq \text{uverts } G$ **and** *fin*: *finite* (*uverts* G)
and *wf*: *wellformed* G
shows $\text{card } (\text{triangle-set } G) \geq 1/6 * \text{card } \{(x,y,z) \in X \times Y \times Z . (\text{triangle-in-graph } x y z G)\}$
 (is - $\geq 1/6 * \text{card } ?TT$)
 <proof>

lemma *card-convert-triangle-rep-bound*:
fixes $G :: \text{ugraph}$ **and** $t :: \text{real}$
assumes $X \subseteq \text{uverts } G$ **and** $Y \subseteq \text{uverts } G$ **and** $Z \subseteq \text{uverts } G$ **and** *fin*: *finite* (*uverts* G)
and *wf*: *wellformed* G
assumes $\text{card } \{(x,y,z) \in X \times Y \times Z . (\text{triangle-in-graph } x y z G)\} \geq t$
shows $\text{card } (\text{triangle-set } G) \geq 1/6 * t$
 <proof>

lemma *edge-density-eq0*:
assumes *all-edges-between* $A B G = \{\}$ **and** $X \subseteq A$ $Y \subseteq B$

shows *edge-density* $X Y G = 0$
 ⟨*proof*⟩

The following is the Triangle Removal Lemma.

theorem *triangle-removal-lemma*:

fixes $\varepsilon :: \text{real}$
assumes *egt*: $\varepsilon > 0$
shows $\exists \delta :: \text{real} > 0. \forall G. \text{card}(\text{uverts } G) > 0 \longrightarrow \text{uwellformed } G \longrightarrow$
 $\text{card}(\text{triangle-set } G) \leq \delta * \text{card}(\text{uverts } G) \wedge 3 \longrightarrow$
 $(\exists G'. \text{triangle-free-graph } G' \wedge \text{uverts } G' = \text{uverts } G \wedge \text{uedges } G' \subseteq \text{uedges}$
 $G \wedge$
 $\text{card}(\text{uedges } G - \text{uedges } G') \leq \varepsilon * (\text{card}(\text{uverts } G))^2)$
(is $\exists \delta :: \text{real} > 0. \forall G. - \longrightarrow - \longrightarrow - \longrightarrow (\exists G \text{new. } ?\Phi G G \text{new}))$
 ⟨*proof*⟩

1.5 Roth's Theorem

We will first need the following corollary of the Triangle Removal Lemma.

See https://en.wikipedia.org/wiki/Ruzsa--Szemerédi_problem. Suggested by Yaël Dillies

corollary *Diamond-free*:

fixes $\varepsilon :: \text{real}$
assumes $0 < \varepsilon$
shows $\exists N > 0. \forall G. \text{card}(\text{uverts } G) > N \longrightarrow \text{uwellformed } G \longrightarrow \text{unique-triangles}$
 $G \longrightarrow$
 $\text{card}(\text{uedges } G) \leq \varepsilon * (\text{card}(\text{uverts } G))^2$
 ⟨*proof*⟩

We are now ready to proceed to the proof of Roth's Theorem for Arithmetic Progressions.

definition *progression3* $:: 'a :: \text{comm-monoid-add} \Rightarrow 'a \Rightarrow 'a \text{ set}$
where *progression3* $k d \equiv \{k, k+d, k+d+d\}$

lemma *p3-int-iff*: *progression3* $(\text{int } k) (\text{int } d) \subseteq \text{int } 'A \iff \text{progression3 } k d \subseteq A$
 ⟨*proof*⟩

We assume that a set of naturals $A \subseteq \{\dots < N\}$ does not have any arithmetic progression. We will then show that A is of cardinality $o(N)$.

lemma *RothArithmeticProgressions-aux*:

fixes $\varepsilon :: \text{real}$
assumes $\varepsilon > 0$
obtains M **where** $\forall N \geq M. \forall A \subseteq \{\dots < N\}. (\nexists k d. d > 0 \wedge \text{progression3 } k d \subseteq A) \longrightarrow \text{card } A < \varepsilon * \text{real } N$
 ⟨*proof*⟩

We finally present the main statement formulated using the upper asymptotic density condition.

theorem *RothArithmeticProgressions:*
 assumes *upper-asymptotic-density* $A > 0$
 shows $\exists k d. d > 0 \wedge \text{progression3 } k d \subseteq A$
 $\langle \text{proof} \rangle$
end