

# The Rogers–Ramanujan Identities

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## Abstract

This entry formalises the Rogers–Ramanujan Identities:

$$\sum_{k=-\infty}^{\infty} \frac{q^{k^2}}{\prod_{j=1}^k (1 - q^j)} = \left( \prod_{n=0}^{\infty} (1 - q^{1+5n})(1 - q^{4+5n}) \right)^{-1}$$
$$\sum_{k=-\infty}^{\infty} \frac{q^{k^2+k}}{\prod_{j=1}^k (1 - q^j)} = \left( \prod_{n=0}^{\infty} (1 - q^{2+5n})(1 - q^{3+5n}) \right)^{-1}$$

The formalisation follows the elegant proof given in Andrews and Eriksson *Integer Partitions*, using the Jacobi triple product.

# 1 The Rogers–Ramanujan identities

```
theory Rogers_Ramanujan
  imports "Theta_Functions.Jacobi_Triple_Product"
begin
```

**Acknowledgement:** I would like to thank George Andrews for giving me a crucial hint about a uniform convergence issue that I struggled with.

```
unbundle qepochhammer_inf_notation
```

First of all, we show two auxiliary results concerned with the (absolute) convergence of two infinite sums that will appear in our proof of the identities.

```
lemma summable_rogers_ramanujan_aux1:
  fixes q :: "'a :: {real_normed_field, banach}" and M :: int
  assumes q: "q ≠ 0" "norm q < 1"
  shows "(λj. norm q powi (j*(5*j+M) div 2)) summable_on UNIV"
proof -
  have *: "(λj. norm q powi (j * (5*j+M) div 2)) summable_on {0..}" for
M :: int
  proof -
    have "summable (λj. norm q powi (int j * (5 * int j + M) div 2))"
    proof (rule summable_comparison_test_bigo)
      have "eventually (λj. int j * (5 * int j + M) div 2 ≥ 0) at_top"
        by real_asymp
      hence "eventually (λj. norm q powi (int j * (5 * int j + M) div
2) =
                                norm q ^ (nat (int j * (5 * int j + M) div
2))) at_top"
        by eventually_elim (auto simp: power_int_def)
      hence "(λj. norm q powi (int j * (5 * int j + M) div 2)) ∈
              Θ(λj. norm q ^ (nat (int j * (5 * int j + M) div 2)))"
        by (rule bigthetaI_cong)
      also have "(λj. norm q ^ (nat (int j * (5 * int j + M) div 2)))
∈ O(λj. (1/2) ^ j)"
        using q by real_asymp
      finally show "(λj. norm q powi (int j * (5 * int j + M) div 2))
∈ O(λj. (1/2) ^ j)" .
    qed auto
    hence "(λj. norm q powi (int j * (5 * int j + M) div 2)) summable_on
UNIV"
      by (rule summable_nonneg_imp_summable_on) auto
    also have "?this ↔ (λj. norm q powi (j * (5*j+M) div 2)) summable_on
{0..}"
      by (rule summable_on_reindex_bij_witness[of _ nat int]) auto
    finally show "(λj. norm q powi (j * (5*j+M) div 2)) summable_on {0..}"
    .
  qed

  have "(λj. norm q powi (j * (5*j-M) div 2)) summable_on {0..}"
```

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    using *[of "-M"] by simp
    also have "?this  $\longleftrightarrow$  ( $\lambda j. \text{norm } q \text{ powi } (j * (5*j+M) \text{ div } 2)$ ) summable_on
    {..0}"
      by (rule summable_on_reindex_bij_witness[of _ uminus uminus]) (auto
    simp: algebra_simps)
    finally have " $(\lambda j. \text{norm } q \text{ powi } (j * (5*j+M) \text{ div } 2))$  summable_on ( $\{..0\}$ 
 $\cup \{0..\}$ )"
      by (intro summable_on_union *)
    also have " $\{..0\} \cup \{0::\text{int}..\} = \text{UNIV}$ "
      by auto
    finally show " $(\lambda j. \text{norm } q \text{ powi } (j*(5*j+M) \text{ div } 2))$  summable_on UNIV"
    .
  qed

```

```

lemma summable_rogers_ramanujan_aux2:
  fixes q :: "'a :: {real_normed_field, banach}" and M :: int
  assumes q: "q  $\neq$  0" "norm q < 1"
  shows "summable ( $\lambda j. \text{norm } (q ^ (j^2 + c * j) / \text{qpochhammer } (\text{int } j) q$ ) q)"
proof (rule summable_comparison_test_bigo)
  have " $(\lambda j. \text{norm } (q ^ (j^2 + c * j) / \text{qpochhammer } (\text{int } j) q)) \in$ 
     $O(\lambda j. \text{norm } q ^ (j^2 + c * j) / (1 - \text{norm } q) ^ j)$ "
  proof (intro bigoI[of _ 1] always_eventually_allI)
    fix j :: nat
    have [simp]: " $\text{qpochhammer } (\text{int } j) q \neq 0$ "
      by (intro qpochhammer_nonneg_nonzero q)
    have "norm (norm (q ^ (j^2 + c * j) / qpochhammer (int j) q q))
    =
      norm q ^ (j^2 + c*j) / norm (qpochhammer (int j) q q)"
      by (simp add: norm_power norm_divide)
    also have "...  $\leq$  norm q ^ (j^2 + c*j) / (1 - norm q) ^ j"
      by (intro divide_left_mono norm_qpochhammer_nonneg_ge mult_pos_pos)
    (use q in auto)
    also have "...  $\leq$  norm (norm q ^ (j^2 + c*j) / (1 - norm q) ^ j)"
      unfolding real_norm_def by linarith
    finally show "norm (norm (q ^ (j^2 + c * j) / qpochhammer (int j) q
    q))
       $\leq$  1 * norm (norm q ^ (j^2 + c * j) / (1 - norm q) ^
    j)"
      by simp
  qed
  also have " $(\lambda j. \text{norm } q ^ (j^2 + c * j) / (1 - \text{norm } q) ^ j) \in O(\lambda j. (1/2) ^ j)$ "
    using q by real_asymp
  finally show " $(\lambda j. \text{norm } (q ^ (j^2 + c * j) / \text{qpochhammer } (\text{int } j) q$ )
  q))  $\in O(\lambda j. (1/2) ^ j)$ " .
qed auto

```

Next, we apply the Jacobi triple product to show that for  $N \in \{0, 1\}$  we

have

$$\sum_{n=-\infty}^{\infty} (-1)^n q^{n(5n+2N+1)/2} = \frac{(q; q)_{\infty}}{\prod_{i \in I} (q^i; q^5)_{\infty}}$$

where  $I = \{1, \dots, 4\} \setminus \{2 - N, 3 + N\}$ .

lemma *rogers\_ramanujan\_aux*:

fixes  $q :: \text{complex}$  and  $N :: \text{nat}$

assumes  $q$ : "*norm*  $q < 1$ " and  $N$ : " $N < 2$ "

shows " $(\lambda n. (-1)^n \text{powi } n * q \text{powi } (n*(5*n+2*N+1) \text{div } 2)) \text{has\_sum}$   
 $(q; q)_{\infty} / ((q^{(1+N)}; q^5)_{\infty} * (q^{(4-N)}; q^5)_{\infty}) \text{UNIV}$ "

proof -

have  $N\_cases$ : " $N \in \{0, 1\}$ "

using  $N$  by auto

have " $(\lambda n. (- (q^{(3+N)})) \text{powi } (n*(n+1) \text{div } 2) * (- (q^{(2-N)})) \text{powi } (n*(n-1) \text{div } 2))$

$\text{has\_sum ramanujan\_theta } (- (q^{(3+N)})) (- (q^{(2-N)})) \text{UNIV}$ "

by (rule *has\_sum\_ramanujan\_theta*)

(use  $q$  in  $\langle \text{simp\_all flip: power\_add add: norm\_power power\_less\_one\_iff} \rangle$ )

also have "*ramanujan\_theta*  $(- (q^{(3+N)})) (- (q^{(2-N)})) =$

$(q^{(3+N)}; q^5)_{\infty} * (q^{(2-N)}; q^5)_{\infty} * (q^5; q^5)_{\infty}$ "

using *ramanujan\_theta\_triple\_product\_complex*[of " $- (q^{(3+N)})$ " " $- (q^{(2-N)})$ "]

$N$

by (*simp flip: power\_add add: norm\_power power\_less\_one\_iff*  $q$ )

also have " $(\lambda n. (- (q^{(3+N)})) \text{powi } (n*(n+1) \text{div } 2) * (- (q^{(2-N)})) \text{powi } (n*(n-1) \text{div } 2)) =$

$(\lambda n. (- 1)^n \text{powi } n * q \text{powi } (n * (5 * n + 2 * \text{int } N + 1) \text{div } 2))$ " (is "*?lhs = ?rhs*")

proof

fix  $n :: \text{int}$

show "*?lhs*  $n = ?rhs$   $n$ "

proof (cases " $q = 0$ ")

case True

have \*: " $n * (n + 1) \text{div } 2 = 0 \iff n \in \{0, -1\}$ "

" $n * (n - 1) \text{div } 2 = 0 \iff n \in \{0, 1\}$ "

by (auto *simp: dvd\_div\_eq\_0\_iff*)

have " $5 * n + 2 * \text{int } N + 1 \neq 0$ "

using  $N$  by *presburger*

hence \*\*: " $n * (5 * n + 2 * \text{int } N + 1) \text{div } 2 = 0 \iff n = 0$ "

using  $N$  by (auto *simp: dvd\_div\_eq\_0\_iff*)

have "*?lhs*  $n = 0 \text{powi } (n * (n + 1) \text{div } 2) * 0 \text{powi } (n * (n - 1) \text{div } 2)$ "

using  $N$  by (*simp add: True power\_0\_left*)

also have "... = (if  $n = 0$  then 1 else 0)"

unfolding *power\_int\_0\_left\_If* \* by auto

also have "... = *?rhs*  $n$ "

unfolding *True power\_int\_0\_left\_If* \*\* by auto

finally show *?thesis* .

next

case [*simp*]: False

```

define m where "m = n*(n-1) div 2"
have "m + (2 * n) div 2 = (n * (n - 1) + 2 * n) div 2"
  unfolding m_def by (subst div_plus_div_distrib_dvd_left [symmetric])
auto
also have "n * (n - 1) + 2 * n = n * (n + 1)"
  by (simp add: algebra_simps)
finally have *: "n * (n + 1) div 2 = m + n"
  by simp

have "(- (q^(3+N))) powi (n*(n+1) div 2) * (- (q^(2-N))) powi (n*(n-1)
div 2) =
  (-1) powi n * (q powi ((3 + N) * (m+n)) * q powi ((2-N) *
m))"
  unfolding * m_def [symmetric]
  by (subst (1 2) power_int_minus_left) (auto simp: power_int_power)
also have "q powi ((3 + N) * (m+n)) * q powi ((2-N) * m) =
  q powi ((3 + N) * (m+n) + (2-N) * m)"
  by (simp add: power_int_add)
also have "(3 + N) * (m+n) + (2-N) * m = N * n + 3 * n + 5 * m"
  using N by (simp add: algebra_simps of_nat_diff)
also have "... = (2 * N * n + 6 * n + 5 * n * (n - 1)) div 2"
  by (simp add: m_def div_mult_swap)
also have "2 * N * n + 6 * n + 5 * n * (n - 1) = n * (5 * n + 2
* int N + 1)"
  by (simp add: algebra_simps)
finally show "?lhs n = ?rhs n" .
qed
qed
finally have "((λn. (-1) powi n * q powi (n*(5*n+2*N+1) div 2)) has_sum
(q^(3+N) ; q^5)_∞ * (q^(2-N) ; q^5)_∞ * (q^5 ; q^5)_∞)
UNIV"
  by simp
also have "(q^(3+N) ; q^5)_∞ * (q^(2-N) ; q^5)_∞ * (q^5 ; q^5)_∞ =
(q ; q)_∞ / ((q^(1+N) ; q^5)_∞ * (q^(4-N) ; q^5)_∞)"
proof -
define A where "A = {2+N, 1-N, 4}"
define B where "B = {...<5} - A"
have q_pow_eq_1_iff: "q ^ k = 1 ↔ k = 0" for k
proof
assume "q ^ k = 1"
hence "norm (q ^ k) = 1"
  by simp
moreover have "k > 0 → norm (q ^ k) < 1"
  using q by (simp add: norm_power power_less_one_iff)
ultimately show "k = 0"
  by auto
qed auto

have "(q^(3+N) ; q^5)_∞ * (q^(2-N) ; q^5)_∞ * (q^5 ; q^5)_∞ =

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      (∏ i∈A. (q ^ (i+1); q ^ 5)∞)"
    using N_cases by (auto simp: eval_nat_numeral mult_ac A_def)
  also have "... = (∏ i∈{..<5} - B. (q ^ (i+1); q ^ 5)∞)"
    using N by (intro prod.cong) (auto simp: A_def B_def)
  also have "... = (∏ i<5. (q ^ (i+1); q ^ 5)∞) / (∏ i∈B. (q ^ (i+1);
q ^ 5)∞)"
    by (subst prod_diff)
      (simp_all add: B_def qpochhammer_inf_zero_iff norm_power power_less_one_iff

          q q_pow_eq_1_iff flip: power_mult power_Suc power_add)
  also have "(∏ i<5. (q ^ (i+1); q ^ 5)∞) = (q; q)∞"
    using prod_qpochhammer_group[of q 5 q] q by simp
  also have "B = {N, 3 - N}"
    unfolding A_def B_def using N_cases by auto
  also have "(∏ i∈{N, 3 - N}. (q ^ (i + 1) ; q ^ 5)∞) = (q^(1+N); q^5)∞
* (q^(4-N); q^5)∞"
    using N by (simp flip: power_Suc flip: Suc_diff_le)
  finally show ?thesis .
qed
finally show ?thesis .
qed

```

```

theorem rogers_ramanujan_complex:
  fixes q :: complex
  assumes "norm q < 1"
  shows "((λj. q ^ (j2) / qpochhammer j q q) has_sum (1 / ((q;q5)∞
* (q4;q5)∞))) UNIV"
  and "((λj. q ^ (j2 + j) / qpochhammer j q q) has_sum (1 / ((q2;q5)∞
* (q3;q5)∞))) UNIV"
proof -
  have "((λj. q ^ (j2) / qpochhammer j q q) has_sum (1 / ((q;q5)∞ *
(q4;q5)∞))) UNIV ∧
  ((λj. q ^ (j2 + j) / qpochhammer j q q) has_sum (1 / ((q2;q5)∞
* (q3;q5)∞))) UNIV"
proof (cases "q = 0")
  case [simp]: True
  show ?thesis
  by (auto intro!: has_sum_finite_neutralI[of "{0}"])
next
  case [simp]: False
  note q = <norm q < 1>
  have q_pow_neq_1: "q ^ n ≠ 1" if "n > 0" for n
    using q_power_neq_1[of q n] that q by simp
  have [simp]: "(q; q)∞ ≠ 0"
    using q qpochhammer_inf_nonzero by blast
  have [simp]: "qpochhammer (int j) q q ≠ 0" for j
proof
  assume "qpochhammer j q q = 0"
  then obtain k where k: "q * q powi k = 1" "k ∈ {0..<int j}"

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```

    using q by (auto simp: qpothhammer_eq_0_iff)
  show False
    using q_pow_neq_1[of "nat (k+1)"] k by (auto simp: nat_add_distrib
power_int_def)
  qed

  define B :: "int  $\Rightarrow$  int  $\Rightarrow$  complex"
    where "B = ( $\lambda$ n k. if  $n \geq 0 \wedge k \geq 0$  then qbinomial q (nat n) (nat
k) else 0)"
  have B_eq: "B (int n) (int k) = qbinomial q n k" for n k
    by (simp add: B_def)
  have [simp]: "B n k = 0" if "n < 0  $\vee$  k < 0  $\vee$  k > n" for n k
    using that by (auto simp: B_def)
  have B_sym: "B n k = B n (n - k)" for n k
    using qbinomial_symmetric[of q "nat k" "nat n"] q
    by (auto simp: B_def nat_diff_distrib)

  have B_rec: "B n k = q powi k * B (n-1) k + B (n-1) (k-1)" if "n  $\neq$ 
0  $\vee$  k  $\neq$  0" for n k :: int
  proof (cases "n > 0  $\wedge$  k > 0")
    case True
      define n' k' where "n' = nat (n-1)" and "k' = nat (k-1)"
      have [simp]: "n = int (Suc n')" "k = int (Suc k')"
        using True unfolding B_def by (auto simp: n'_def k'_def)
      show ?thesis
        by (auto simp: B_def nat_add_distrib power_int_def qbinomial_Suc_Suc)
      qed (use that in <auto simp: B_def qbinomial_0_middle>)

    case False
      have B_rec': "B n k = B (n-1) k + q powi (n-k) * B (n-1) (k-1)" if
"n  $\neq$  0  $\vee$  k  $\neq$  0" for n k :: int
      proof (cases "n > 0  $\wedge$  k > 0")
        case True
          define n' k' where "n' = nat (n-1)" and "k' = nat (k-1)"
          have [simp]: "n = int (Suc n')" "k = int (Suc k')"
            using True unfolding B_def by (auto simp: n'_def k'_def)
          show ?thesis using q
            by (auto simp: B_def nat_add_distrib power_int_def qbinomial_Suc_Suc'
nat_diff_distrib)
          qed (use that in <auto simp: B_def qbinomial_0_middle>)

        case False
          have B_rec3: "B n k * (1 - q ^ (n - k)) = (1 - q ^ n) * B (n - 1)
k" for n k
          proof (cases "n > 0  $\wedge$  k > 0  $\wedge$  k < n")
            case True
              thus ?thesis
                using qbinomial_rec2[of q "nat n" "nat k"] q
                by (auto simp: B_def q_pow_neq_1 nat_diff_distrib)
            case False
              qed (auto simp: B_def)
          end
        end
      end
    end
  end

```

```

define e :: "int ⇒ int ⇒ int" where "e = (λN j. j*(5 * j + 1 -
4 * N) div 2)"
define S where "S = (λN n. ∑ j ≤ n. q^(j^2 + nat N * j) * B n j)"
define T where "T = (λN n. ∑ ∞ j. (-1) powi j * q powi (e N j) *
B (2 * int n + N) (int n + 2*j))"

have e0_conv_e1: "e 0 j = e 1 j + 2 * j" for j
proof -
  have "e 0 j = j * (5 * j + 1) div 2"
    by (simp add: e_def)
  also have "j * (5 * j + 1) = j * (5 * j - 3) + 4 * j"
    by (simp add: algebra_simps)
  also have "... div 2 = e 1 j + 2 * j"
    by (subst div_plus_div_distrib_dvd_left) (auto simp: e_def)
  finally show ?thesis .
qed

have has_sum_aux:
  "(λj. (-1) powi j * q powi (e a j) * B (2 * int n + a+b) (int (n+b)
+ 2*j)) summable_on UNIV"
  (is "?f summable_on _") for a b n
proof -
  have "finite {j. 2 * j ∈ {-int (n+b)..int n + a}}"
  proof (rule finite_subset)
    show "{j. 2 * j ∈ {-int (n+b)..int n + a}} ⊆
      {-(n+b+1) div 2..(int n + a + 1) div 2}"
      by (auto simp: minus_le_iff)
  qed auto
  hence "?f summable_on {j. 2 * j ∈ {-int (n+b)..int n + a}}"
    by (rule summable_on_finite)
  also have "?this ⟷ ?f summable_on UNIV"
    by (rule summable_on_cong_neutral) auto
  finally show ?thesis .
qed

have has_sum_T: "((λj. (-1) powi j * q powi (e N j) * B (2*n+N) (n+2*j))
has_sum T N n) UNIV"
  for n N using has_sum_infsum[OF has_sum_aux[of N n 0]] by (simp
add: T_def)

have has_sum_T0: "((λj. (-1) powi j * q powi (e 0 j) * B (2*n+1) (n+1+2*j))
has_sum T 0 n) UNIV"
  for n
proof -
  define T' where
    "T' = (λn. ∑ ∞ j. (-1) powi j * q powi (e 0 j) * B (2 * int n
+ 1) (int n+1+2*j))"
  have 1: "((λj::int. (-1) powi j * q powi e 0 j * B (2*n+1) (n+1+2*j))
has_sum T' n) UNIV"

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    using has_sum_infsum[OF has_sum_aux[of 0 n 1]] unfolding T'_def
  by (simp add: algebra_simps)
    also have "(λj::int. (-1) powi j * q powi e 0 j * B (2*n+1) (n+1+2*j))
=
    (λj::int. (-1) powi j * q powi e 0 j *
      (B (2*n) (n+2*j) + q powi (n+1+2*j) * B (2*n) (n+1+2*j)))"
    by (subst B_rec) auto
  finally have 2: "... has_sum (T' n) UNIV" .

  have 3: "((λj::int. (-1) powi j * q powi (e 0 j + 2 * j) * q powi
(n+1) * B (2*n) (n+2*j+1))
    has_sum (T' n - T 0 n)) UNIV"
    using has_sum_diff[OF 2 has_sum_T[of 0 n]] by (simp add: algebra_simps
power_int_add)
    also have "?this ⟷ ((λj::int. -((-1) powi j * q powi (e 0 (-j-1)
- 2*j-2) * q powi (n+1) * B (2*n) (n-2*j-1)))
    has_sum (T' n - T 0 n)) UNIV"
    by (intro has_sum_reindex_bij_witness[of _ "λj. -j-1" "λj. -j-1"])
      (auto simp: algebra_simps power_int_diff power_int_minus power_int_minus_left)
    also have "(λj::int. -((-1) powi j * q powi (e 0 (-j-1) - 2*j-2)
* q powi (n+1) * B (2*n) (n-2*j-1))) =
      (λj::int. -((-1) powi j * q powi (e 0 j + 2*j) * q powi
(n+1) * B (2*n) (n+2*j+1)))"
      (is "?lhs = ?rhs")
    proof
      fix j :: int
      have "?lhs j = - ((- 1) powi j * q powi (e 0 (-j-1) - (4*j+2)
+ 2 * j) *
      q powi int (n + 1) * B (2*n) (n+2*j+1))"
      by (subst B_sym) (simp_all add: algebra_simps)
      also have "e 0 (-j-1) - (4*j+2) = ((j+1)*(4+5*j) - 2*(4*j+2))
div 2"
      unfolding e_def by (subst div_diff) (auto simp: algebra_simps)
      also have "(j+1)*(4+5*j) - 2*(4*j+2) = j * (5*j + 1)"
      by (simp add: algebra_simps)
      also have "... div 2 = e 0 j"
      by (simp add: e_def algebra_simps)
      finally show "?lhs j = ?rhs j" .
    qed
  finally have "((λj::int. (-1) powi j * q powi (e 0 j + 2*j) * q powi
(n+1) * B (2*n) (n+2*j+1))
    has_sum (T 0 n - T' n)) UNIV"
    by (subst (asm) has_sum_uminus) simp_all
  with 3 have "T' n - T 0 n = T 0 n - T' n"
    using has_sum_unique by blast
  hence "T 0 n = T' n"
    by simp
  with 1 show ?thesis
    by (simp add: algebra_simps)

```

qed

The initial conditions  $S_i(0) = T_i(0) = 1$  are easily determined.

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have [simp]: "S N 0 = 1" for N
  by (simp_all add: S_def B_def)

have [simp]: "T N 0 = 1" if N: "N ∈ {0,1}" for N
proof -
  have "((λj. (-1) powi j * q powi e N j * B N (2 * j)) has_sum 1)
{0}"
    by (intro has_sum_finiteI) (use N in <auto simp: e_def B_def>)
  also have "?thesis ↔ ((λj. (-1) powi j * q powi e N j * B N (2
* j)) has_sum 1) UNIV"
    using N by (intro has_sum_cong_neutral) (auto simp: B_def)
  finally show ?thesis
    by (simp add: has_sum_iff T_def)
qed

```

Next, we prove that the  $S_i$  satisfy the recurrences  $S_0(n+1) = S_0(n) + q^{n+1}S_1(n)$  and  $S_1(n+1) = q^{n+1}S_0(n+1) + (1 - q^{n+1})S_1(n)$ . This requires just a few manipulations of finite sums involving the two recurrences for the  $q$ -binomial coefficients.

```

have S0_rec: "S 0 (Suc n) = S 0 n + q ^ Suc n * S 1 n" for n
proof -
  define n' where "n' = Suc n"
  have n': "n' > 0" and [simp]: "n' - 1 = n" "int n' - 1 = int n"
    by (simp_all add: n'_def)

  have "S 0 n' = (∑ j ≤ n'. q ^ j^2 * B (int n') (int j))"
    by (simp add: S_def)
  also have "... = (∑ j ≤ n'. q^(j^2) * B (int n) (int j)) +
(∑ j ≤ n'. q powi (int j ^ 2 - int j) * q ^ n' *
B (int n) (int j - 1))"
    by (subst B_rec')
    (use n' in <simp_all add: algebra_simps power_add sum.distrib
power_int_diff
power_int_add flip: of_nat_power>)
  also have "(∑ j ≤ n'. q^(j^2) * B (int n) (int j)) = S 0 n"
    unfolding S_def by (intro sum_mono_neutral_cong_right) (auto simp:
n'_def)
  also have "(∑ j ≤ n'. q powi (int j ^ 2 - int j) * q ^ n' * B (int
n) (int j - 1)) =
q^n' * (∑ j ≤ n'. q powi (int j ^ 2 - int j) * B (int
n) (int j - 1))"
    by (simp add: sum_distrib_left sum_distrib_right mult_ac)
  also have "(∑ j ≤ n'. q powi (int j ^ 2 - int j) * B (int n) (int
j - 1)) =
(∑ j=1..n'. q powi (int j ^ 2 - int j) * B (int n) (int
j - 1))"

```

```

    by (intro sum.mono_neutral_right) auto
    also have "... = (∑ j ≤ n. q powi (int j ^ 2 + int j) * B (int n)
(int j))"
    by (intro sum.reindex_bij_witness[of _ "λj. j+1" "λj. j-1"])
      (auto simp: n'_def of_nat_diff power2_eq_square algebra_simps)
    also have "... = (∑ j ≤ n. q ^ (j^2+j) * B (int n) (int j))"
    by (rule sum.cong) (auto simp: power_int_def nat_add_distrib nat_power_eq)
    also have "q ^ n' * ... = q ^ n' * S 1 n"
    by (simp add: S_def)
    finally show ?thesis
    unfolding n'_def .
qed

```

```

have S1_rec: "S 1 (Suc n) = q ^ (Suc n) * S 0 (Suc n) + (1 - q ^ Suc
n) * S 1 n" for n
proof -
  define n' where "n' = Suc n"
  have n': "n' > 0" and [simp]: "n' - 1 = n" "int n' - 1 = int n"
  by (simp_all add: n'_def)
  have "S 1 n' - q^n' * S 0 n' = (∑ j ≤ n'. q ^ (j^2 + j) * (B (int
n') (int j) * (1 - q ^ (n' - j))))"
  by (simp add: S_def sum_distrib_left power_int_add ring_distrib
power_diff
      power_add mult_ac flip: sum_subtractf)
  also have "... = (∑ j ≤ n. q ^ (j^2 + j) * (B (int n') (int j) * (1
- q ^ (n' - j))))"
  by (rule sum.mono_neutral_right) (auto simp: n'_def)
  also have "... = (∑ j ≤ n. q ^ (j^2 + j) * B (int n' - 1) (int j)
* (1 - q ^ n'))"
  by (subst B_rec3) (auto simp: mult_ac)
  also have "... = (1 - q ^ n') * S 1 n"
  by (simp add: S_def sum_distrib_left sum_distrib_right mult_ac)
  finally show ?thesis
  unfolding n'_def by (simp add: algebra_simps)
qed

```

Next, we show that the  $T_i$  satisfy equivalent recurrence relations. This again consists of some conceptually simple manipulations of sums. The sums are again finite in the sense that all but finitely many summands are 0, but it is a bit easier to simply sum over all integers and not worry about the precise summation range.

```

have T0_rec: "T 0 (Suc n) = T 0 n + q ^ Suc n * T 1 n" for n
proof -
  have "((λj::int. (-1) powi j * q powi e 0 j * B (2*n+2) (n+1+2*j)
-
      (-1) powi j * q powi e 0 j * B (2*n+1) (n+1+2*j))
      has_sum (T 0 (Suc n) - T 0 n)) UNIV"
  using has_sum_T[of 0 "Suc n"] has_sum_T0[of n]
  by (intro has_sum_diff) (simp_all add: algebra_simps)

```

```

    hence "((λj::int. (-1) powi j * q powi e 0 j * (B (2*n+2) (n+1+2*j)
- B (2*n+1) (n+1+2*j))))
      has_sum (T 0 (Suc n) - T 0 n) UNIV"
    by (simp add: algebra_simps)
    also have "(λj::int. B (2*n+2) (n+1+2*j) - B (2*n+1) (n+1+2*j))
=
      (λj::int. B (2*n+1) (n+2*j) * q powi (n+1-2*j))"
    by (subst B_rec') (simp_all add: algebra_simps)
    finally have "(λj::int. (-1) powi j * q powi e 0 j * B (2*n+1) (n+2*j)
* q powi (n+1-2*j))
      has_sum T 0 (Suc n) - T 0 n UNIV"
    by (simp only: mult.assoc)
    also have "(λj::int. (-1) powi j * q powi e 0 j * B (2*n+1) (n+2*j)
* q powi (n+1-2*j)) =
      (λj::int. q powi (n+1) * ((-1) powi j * q powi (e 1 j)
* B (2*n+1) (n+2*j))))"
    (is "?lhs = ?rhs")
  proof
    fix j :: int
    have "?lhs j = q powi (n+1) * ((-1) powi j * q powi (e 0 j - 2
* j) * B (2*n+1) (n+2*j))"
      by (simp add: power_int_add power_int_diff field_simps)
    also have "e 0 j - 2 * j = e 1 j"
      by (simp add: e0_conv_e1)
    finally show "?lhs j = ?rhs j" .
  qed
  finally have "(λj::int. q powi (n+1) * ((-1) powi j * q powi (e
1 j) * B (2*n+1) (n+2*j)))
    has_sum (T 0 (Suc n) - T 0 n) UNIV" .
  moreover have "(λj::int. q powi (n+1) * ((-1) powi j * q powi
(e 1 j) * B (2*n+1) (n+2*j))) has_sum
    (q powi (n+1) * T 1 n) UNIV"
    using has_sum_T[of 1 n] by (intro has_sum_cmult_right has_sum_T)
(simp_all add: add_ac)
  ultimately have "T 0 (Suc n) - T 0 n = q powi (n+1) * T 1 n"
    using has_sum_unique by blast
  thus ?thesis
    by (simp add: algebra_simps power_int_add)
qed

  have T1_rec: "T 1 (Suc n) = q ^ Suc n * T 0 (Suc n) + (1 - q ^ Suc
n) * T 1 n" for n
  proof -
    have "(λj::int. (-1) powi j * q powi e 1 j * B (2*n+3) (n+1+2*j)
-
      q ^ Suc n * ((-1) powi j * q powi e 0 j * B (2*n+2)
(n+1+2*j)))
      has_sum (T 1 (Suc n) - q ^ Suc n * T 0 (Suc n))
UNIV"

```

```

    using has_sum_T[of 1 "Suc n"] has_sum_T[of 0 "Suc n"]
    by (intro has_sum_diff has_sum_cmult_right) (simp_all add: algebra_simps)
    also have "(λj::int. (-1) powi j * q powi e 1 j * B (2*n+3) (n+1+2*j)
-
      q ^ Suc n * ((-1) powi j * q powi e 0 j * B (2*n+2)
(n+1+2*j))) =
      (λj::int. (-1) powi j * q powi e 1 j *
      (B (2*n+1) (n+2*j) + q powi (n+2-2*j) * B (2*n+1)
(n+2*j-1)))"
      (is "?lhs = ?rhs")
    proof
      fix j :: int
      have "?lhs j = (-1) powi j * (q powi e 1 j * B (2*n+3) (n+1+2*j)
-
        q powi (e 0 j - 2*j) * q powi (n+1+2*j) * B
(2*n+2) (n+1+2*j)))"
        by (simp add: algebra_simps power_int_add power_int_diff)
      also have "e 0 j - 2 * j = e 1 j"
        by (simp add: e0_conv_e1)
      also have "(-1) powi j * (q powi e 1 j * B (2*n+3) (n+1+2*j) -
q powi e 1 j * q powi (n+1+2*j) * B (2*n+2) (n+1+2*j))
=
      (-1) powi j * q powi e 1 j *
      (B (2*n+3) (n+1+2*j) - q powi (n+1+2*j) * B (2*n+2)
(n+1+2*j))"
        by (simp add: algebra_simps)
      also have "... = (-1) powi j * q powi e 1 j * B (2*n+2) (n+2*j)"
        by (subst B_rec) (simp_all add: algebra_simps)
      also have "... = (-1) powi j * q powi e 1 j *
      (B (2*n+1) (n+2*j) + q powi (n+2-2*j) * B (2*n+1)
(n+2*j-1))"
        by (subst B_rec') (simp_all add: algebra_simps)
      finally show "?lhs j = ?rhs j" .
    qed
    finally have 1: "(λj::int. (-1) powi j * q powi e 1 j *
      (B (2*n+1) (n+2*j) + q powi (n+2-2*j) * B (2*n+1)
(n+2*j-1)))
      has_sum (T 1 (Suc n) - q ^ Suc n * T 0 (Suc n))
UNIV" .

    have 2: "(λj::int. (-1) powi j * q powi e 1 j * B (2*n+1) (n+2*j))
has_sum T 1 n) UNIV"
      using has_sum_T[of 1 n] by (simp add: add_ac)
    have "(λj::int. q^n * ((-1) powi j * q powi (e 1 j - 2*j + 2) *
B (2*n+1) (n+2*j-1)))
      has_sum (T 1 (Suc n) - q ^ Suc n * T 0 (Suc n)
- T 1 n) UNIV"
      using has_sum_diff[OF 1 2]
      by (simp add: algebra_simps power_int_add power_int_diff power2_eq_square)

```

```

    also have "?this  $\longleftrightarrow$  (( $\lambda j :: \text{int. } -(q^{(\text{Suc } n)} * ((-1)^{\text{powi } j} * q^{\text{powi } (e 1 (1-j) + 2*j - 1)} * B (2*n+1) (n-2*j+1))))$ )
      has_sum (T 1 (Suc n) - q ^ Suc n * T 0 (Suc
n) - T 1 n)) UNIV"
    by (intro has_sum_reindex_bij_witness[of _ " $\lambda j. 1 - j$ " " $\lambda j. 1 - j$ "])
      (auto simp: power_int_diff power_int_add power_int_minus_left
algebra_simps power2_eq_square)
    also have " $(\lambda j. e 1 (1 - j) + 2 * j - 1) = (\lambda j. e 1 j)$ "
    proof
      fix j :: int
      have " $e 1 (1 - j) + 2 * j - 1 = ((1 - j) * (2 - 5 * j) + (4 * j - 2)) \text{ div } 2$ "
      unfolding e_def by (subst div_plus_div_distrib_dvd_left) auto
      also have " $(1 - j) * (2 - 5 * j) + (4 * j - 2) = j * (5 * j - 3)$ "
      by (simp add: algebra_simps)
      also have "... div 2 = e 1 j"
      by (simp add: algebra_simps e_def)
      finally show " $e 1 (1 - j) + 2 * j - 1 = e 1 j$ ".
    qed
    also have " $(\lambda j :: \text{int. } B (\text{int } (2 * n + 1)) (\text{int } n - 2 * j + 1)) =$ 
      ( $\lambda j. B (\text{int } (2*n+1)) (\text{int } n + 2*j)$ )"
      by (subst B_sym) (auto simp: algebra_simps)
    finally have " $((\lambda j :: \text{int. } (q^{(\text{Suc } n)} * ((-1)^{\text{powi } j} * q^{\text{powi } e 1 j} * B (2*n+1) (n+2*j))))$ )
      has_sum -(T 1 (Suc n) - q ^ Suc n * T 0 (Suc n)
- T 1 n)) UNIV"
      by (simp only: has_sum_uminus)

    moreover have " $((\lambda j :: \text{int. } (q^{(\text{Suc } n)} * ((-1)^{\text{powi } j} * q^{\text{powi } e 1 j} * B (2*n+1) (n+2*j))))$ )
      has_sum (q^{(\text{Suc } n)} * T 1 n)) UNIV"
      using has_sum_T[of 1 n] by (intro has_sum_cmult_right) (simp_all
add: add_ac)
    ultimately have " $-(T 1 (\text{Suc } n) - q ^ \text{Suc } n * T 0 (\text{Suc } n) - T 1 n)$ 
=  $q^{(\text{Suc } n)} * T 1 n$ "
      using has_sum_unique by blast
    thus ?thesis
      by (simp add: algebra_simps)
  qed

  have S0_eq: "S 0 = T 0" and S1_eq: "S 1 = T 1"
  proof -
    have "S 0 n = T 0 n  $\wedge$  S 1 n = T 1 n" for n
      by (induction n) (simp_all add: S0_rec T0_rec S1_rec T1_rec)
    thus "S 0 = T 0" "S 1 = T 1"
  
```

```

    by blast+
qed

have B_lim1: "(λn. B (f n) (int m)) → 1 / qpochhammer m q q"
  if "filterlim f at_top at_top" for f :: "nat ⇒ int" and m :: nat
proof -
  have f: "eventually (λn. f n ≥ 0) at_top"
    using that by (simp add: filterlim_at_top)
  have "(λn. qbinomial q n m) → 1 / qpochhammer m q q"
    by (rule tendsto_qbinomial1) (use q in auto)
  moreover have "filterlim (λn. nat (f n)) at_top at_top"
    by (rule filterlim_compose[OF _ that]) real_asymp
  ultimately have "(λn. qbinomial q (nat (f n)) m) → 1 / qpochhammer
m q q"
    by (rule filterlim_compose)
  also have "eventually (λn. qbinomial q (nat (f n)) m = B (f n) (int
m)) at_top"
    using f by eventually_elim (auto simp: B_def)
  hence "(λn. qbinomial q (nat (f n)) m) → 1 / qpochhammer m
q q ↔ ?thesis"
    by (intro filterlim_cong) auto
  finally show ?thesis .
qed

have B_lim2: "(λn. B (f n) (g n)) → 1 / (q; q)∞"
  if "filterlim (λn. f n - g n) at_top at_top" "filterlim g at_top
at_top"
  for f g :: "nat ⇒ int"
proof -
  have g_pos: "eventually (λn. g n > 0) at_top"
    using that(2) eventually_compose_filterlim eventually_gt_at_top
by blast
  have "eventually (λn. f n - g n > 0) at_top"
    using that(1) eventually_compose_filterlim eventually_gt_at_top
by blast
  hence fg: "eventually (λn. f n ≥ g n) at_top"
    by eventually_elim auto

  have "(λn. qbinomial q (nat (f n)) (nat (g n))) → 1 / (q; q)∞"
proof (rule tendsto_qbinomial2)
  show "filterlim (λn. nat (g n)) at_top at_top"
    by (rule filterlim_compose[OF _ that(2)]) real_asymp
  have "filterlim (λn. nat (f n - g n)) at_top at_top"
    by (rule filterlim_compose[OF _ that(1)]) real_asymp
  also have "eventually (λn. nat (f n - g n) = nat (f n) - nat (g
n)) at_top"
    using fg g_pos by eventually_elim (auto simp: nat_diff_distrib)
  hence "filterlim (λn. nat (f n - g n)) at_top at_top ↔
filterlim (λn. nat (f n) - nat (g n)) at_top at_top"

```

```

    by (intro filterlim_cong) auto
    finally show "filterlim (λn. nat (f n) - nat (g n)) at_top at_top"
.
    qed (use q in auto)
    also have "eventually (λn. qbinomial q (nat (f n)) (nat (g n))) =
B (f n) (g n) at_top"
    using g_pos fg by eventually_elim (auto simp: B_def)
    hence "(λn. qbinomial q (nat (f n)) (nat (g n))) → 1 / (q;q)∞
←→
    (λn. B (f n) (g n)) → 1 / (q;q)∞"
    by (intro filterlim_cong) auto
    finally show ?thesis .
qed

```

The  $q$ -binomial coefficient is bounded uniformly for all  $n, k$  (with a fixed  $q$ ):

```

define c where "c = (-norm q; norm q)∞ / (norm q ; norm q)∞"
have B_bound: "norm (B n k) ≤ c" for n k
proof (cases "0 ≤ k ∧ k ≤ n")
case True
thus ?thesis
using norm_qbinomial_le_qpochhammer_inf[of q "nat n" "nat k"]
q
by (auto simp: B_def c_def)
next
case False
hence "norm (B n k) = 0"
by auto
also have "0 ≤ c"
unfolding c_def using q by (auto intro!: divide_nonneg_nonneg
qpochhammer_inf_nonneg)
finally show ?thesis .
qed

have uniform_limit1:
"uniform_limit UNIV (λX n. ∑ j∈X. q ^ (j2 + N * j) * B (int n)
(int j)) (S N) finite_sets_at_top"
for N
proof (rule Weierstrass_m_test_general')
fix n :: nat
have "((λj. q ^ (j2 + N * j) * B (int n) (int j)) has_sum S N n)
{..n}"
by (intro has_sum_finiteI) (auto simp: S_def)
also have "?this ↔ ((λj. q ^ (j2 + N * j) * B (int n) (int j))
has_sum S N n) UNIV"
by (intro has_sum_cong_neutral) auto
finally show "((λj. q ^ (j2 + N * j) * B (int n) (int j)) has_sum
S N n) UNIV" .
next
fix j n :: nat

```



```

    have "norm (q ^ (j^2 + N*j) * B n j) = norm q ^ (j^2 + N*j) * norm
(B n j)"
      by (simp add: norm_power norm_mult)
    also have "... ≤ norm q ^ (j^2+N*j) * c"
      by (intro mult_left_mono norm_qbinomial_le B_bound) auto
    finally show "norm (q ^ (j^2+N*j) * B (int n) (int j)) ≤ c * norm
q ^ (j^2+N*j)"
      by (simp add: algebra_simps)
  next
    show "(λj. c * norm q ^ (j^2+N*j)) summable_on UNIV"
    proof (intro summable_on_cmult_right summable_nonneg_imp_summable_on)
      show "summable (λj. norm q ^ (j^2 + N*j))"
      proof (rule summable_comparison_test_bigo)
        show "(λj. norm q ^ (j^2+N*j)) ∈ O(λn. (1/2) ^ n)"
          using q by real_asymp
      qed auto
    qed auto
  qed
qed

  have uniform_limit2:
    "uniform_limit UNIV (λJ n. ∑ j∈J. (-1) powi j * q powi e N j *
B (2*n+N) (n+2*j)) (T N)
    finite_sets_at_top" for N
  proof (rule Weierstrass_m_test_general'[OF _ has_sum_T])
    fix j :: int and n :: nat
    have "norm ((-1) powi j * q powi e N j * B (2*int n+N) (int n +
2*j)) =
      norm q powi e N j * norm (B (2*int n+N) (int n + 2*j))"
      by (simp add: norm_mult norm_power_int)
    also have "... ≤ norm q powi e N j * c"
      by (intro mult_left_mono B_bound) auto
    finally show "norm ((-1) powi j * q powi e N j * B (2*int n+N) (int
n + 2*j)) ≤
      c * norm q powi e N j"
      by (simp add: mult_ac)
  next
    have "(λj. norm q powi e N j) summable_on UNIV"
      using summable_rogers_ramanujan_aux1[of q "1-4*N"] q by (simp
add: e_def algebra_simps)
    thus "(λj. c * norm q powi e N j) summable_on UNIV"
      by (rule summable_on_cmult_right)
  qed

  have tendsto_S: "S N → (∑ ∞j. q ^ (j^2 + N*j) / qepochhammer j
q q)" for N
  proof (rule swap_uniform_limit'[OF _ _ uniform_limit1]; (intro always_eventually
allI)?)
  next
    fix X :: "nat set"

```

```

have "(λn. ∑ j∈X. q ^ (j^2+N*j) * B (int n) (int j)) →
      (∑ j∈X. q ^ (j^2+N*j) * (1 / qpochhammer j q q))"
  by (intro tendsto_intros B_lim1)
thus "(λn. ∑ j∈X. q ^ (j^2+N*j) * B (int n) (int j)) →
      (∑ j∈X. q ^ (j^2+N*j) / qpochhammer j q q)"
  by simp
next
have "summable (λj. norm (q ^ (j^2+N*j) / qpochhammer (int j) q q))"
  using summable_rogers_ramanujan_aux2[of q N] q by simp
hence "(λj. q ^ (j^2+N*j) / qpochhammer j q q) has_sum
      (∑ ∞ j. q ^ (j^2+N*j) / qpochhammer j q q) UNIV"
  by (intro has_sum_infsup norm_summable_imp_summable_on)
thus "(sum (λj. q ^ (j^2+N*j) / qpochhammer (int j) q q) →
      (∑ ∞ j. q ^ (j^2+N*j) / qpochhammer j q q)) finite_sets_at_top"
  by (simp add: has_sum_def)
qed auto

have tendsto_T0: "T 0 → (1 / ((q; q^5)∞ * (q^4; q^5)∞))"
proof (rule swap_uniform_limit'[OF _ _ uniform_limit2]; (intro always_eventually
allI?))
  fix X :: "int set"
  show "(λn. ∑ j∈X. (-1) powi j * q powi e 0 j * B (2*n+0) (n+2*j))
→
      (∑ j∈X. (-1) powi j * q powi e 0 j * (1 / (q; q)∞))"
  by (intro tendsto_intros B_lim2) real_asymp+
next
have "(λx. (-1) powi x * q powi e 0 x) has_sum
      ((q; q)∞ / ((q; q^5)∞ * (q^4; q^5)∞)) UNIV"
  unfolding e_def using rogers_ramanujan_aux[of q 0] q by simp
hence "(λj::int. (-1) powi j * q powi e 0 j * (1 / (q; q)∞))
      has_sum ((q; q)∞ / ((q; q^5)∞ * (q^4; q^5)∞) * (1 /
(q; q)∞)) UNIV"
  by (rule has_sum_cmult_left)
also have "(q; q)∞ / ((q; q^5)∞ * (q^4; q^5)∞) * (1 / (q; q)∞)
=
      1 / ((q; q^5)∞ * (q^4; q^5)∞)"
  by simp
finally show "(sum (λj. (-1) powi j * q powi e 0 j * (1 / (q; q)∞))
→
      (1 / ((q; q^5)∞ * (q^4; q^5)∞))) finite_sets_at_top"
  by (simp add: has_sum_def)
qed auto

have tendsto_T1: "T 1 → (1 / ((q^2; q^5)∞ * (q^3; q^5)∞))"
proof (rule swap_uniform_limit'[OF _ _ uniform_limit2]; (intro always_eventually
allI?))
  fix X :: "int set"
  show "(λn. ∑ j∈X. (-1) powi j * q powi e 1 j * B (2*n+1) (n+2*j))

```

```

→
      (∑ j∈X. (-1) powi j * q powi e 1 j * (1 / (q; q)∞))"
    by (intro tendsto_intros B_lim2) real_asymp+
next
  have "((λj. (-1) powi j * q powi ((j*(5*j+3)) div 2)) has_sum
    ((q; q)∞ / ((q^2; q^5)∞ * (q^3; q^5)∞))) UNIV"
    unfolding e_def using rogers_ramanujan_aux[of q 1] q
    by (simp add: algebra_simps eval_nat_numeral)
  also have "?this ↔ ((λx. (-1) powi x * q powi e 1 x) has_sum
    ((q; q)∞ / ((q^2; q^5)∞ * (q^3; q^5)∞)))
UNIV"
    proof (intro has_sum_reindex_bij_witness[of _ uminus uminus] refl)
      fix j :: int
      have "(-1) powi j * q powi (j*(5*j+3) div 2) = (-1) powi (-j)
* q powi (j*(5*j+3) div 2)"
        by (auto simp: power_int_minus field_simps)
      moreover have "(j * (5 * j + 3)) div 2 = e 1 (-j)"
        unfolding e_def by (rule arg_cong[of _ _ "λn. n div 2"]) (simp_all
add: algebra_simps)
      ultimately show "(-1) powi -j * q powi e 1 (-j) = (-1) powi j
* q powi (j*(5*j+3) div 2)"
        by simp
      qed auto
    finally have "((λx. (-1) powi x * q powi e 1 x) has_sum
      ((q; q)∞ / ((q^2; q^5)∞ * (q^3; q^5)∞))) UNIV"
    .
  hence "((λj::int. (-1) powi j * q powi e 1 j * (1 / (q ; q)∞))
    has_sum ((q; q)∞ / ((q^2; q^5)∞ * (q^3; q^5)∞) * (1
/ (q ; q)∞))) UNIV"
    by (rule has_sum_cmult_left)
  also have "(q; q)∞ / ((q^2; q^5)∞ * (q^3; q^5)∞) * (1 / (q ; q)∞)
=
    1 / ((q^2; q^5)∞ * (q^3; q^5)∞)"
    by simp
  finally show "(sum (λj. (-1) powi j * q powi e 1 j * (1 / (q ; q)∞))
→
    (1 / ((q^2; q^5)∞ * (q^3; q^5)∞))) finite_sets_at_top"
    by (simp add: has_sum_def)
  qed auto

```

Now we need only combine everything we have shown and we're done.

```

have "(∑ ∞j. q ^ (j^2) / qpochhammer j q q) = 1 / ((q;q^5)∞ * (q^4;q^5)∞)"
  using tendsto_S[of 0] tendsto_TO LIMSEQ_unique
  unfolding of_nat_0 mult_0 add_0_right S0_eq S1_eq
  by blast
moreover have "(λj. q ^ (j^2) / qpochhammer j q q) summable_on UNIV"
  by (rule norm_summable_imp_summable_on)
  (use summable_rogers_ramanujan_aux2[of q 0] q in simp_all)

```

```

ultimately have th1:
  "((λj. q ^ (j^2) / qpochhammer j q q) has_sum (1 / ((q;q^5)_∞
* (q^4;q^5)_∞))) UNIV"
  by (simp add: has_sum_iff)

  have "(∑_∞ j. q ^ (j^2 + j) / qpochhammer j q q) = 1 / ((q^2 ; q^5)_∞
* (q^3 ; q^5)_∞)"
    using tendsto_S[of 1] tendsto_T1 LIMSEQ_unique
    unfolding of_nat_1 mult_1 S0_eq S1_eq
    by blast
  moreover have "(λj. q ^ (j^2 + j) / qpochhammer j q q) summable_on
UNIV"
    by (rule norm_summable_imp_summable_on)
      (use summable_rogers_ramanujan_aux2[of q 1] q in simp_all)
  ultimately have th2:
    "((λj. q ^ (j^2 + j) / qpochhammer j q q) has_sum (1 / ((q^2;q^5)_∞
* (q^3;q^5)_∞))) UNIV"
    by (simp add: has_sum_iff)

  from th1 and th2 show ?thesis
    by blast
qed
thus "((λj. q ^ (j^2) / qpochhammer j q q) has_sum (1 / ((q;q^5)_∞ *
(q^4;q^5)_∞))) UNIV"
  and "((λj. q ^ (j^2 + j) / qpochhammer j q q) has_sum (1 / ((q^2;q^5)_∞
* (q^3;q^5)_∞))) UNIV"
  by blast+
qed

lemma rogers_ramanujan_real:
  fixes q :: real
  assumes "|q| < 1"
  shows "((λj. q ^ (j^2) / qpochhammer j q q) has_sum (1 / ((q;q^5)_∞
* (q^4;q^5)_∞))) UNIV"
  and "((λj. q ^ (j^2 + j) / qpochhammer j q q) has_sum (1 / ((q^2;q^5)_∞
* (q^3;q^5)_∞))) UNIV"
proof -
  define q' where "q' = complex_of_real q"
  have [simp]: "norm q' < 1"
    using assms by (simp add: q'_def)

  have "((λj. q' ^ (j^2) / qpochhammer j q' q') has_sum (1 / ((q';q'^5)_∞
* (q'^4;q'^5)_∞))) UNIV"
    by (rule rogers_ramanujan_complex) auto
  also have "(λj. q' ^ (j^2) / qpochhammer j q' q') = (λj. of_real (q ^
(j^2) / qpochhammer j q q))"
    by (simp add: q'_def qpochhammer_of_real)
  also have "(1 / ((q';q'^5)_∞ * (q'^4;q'^5)_∞)) = of_real (1 / ((q;q^5)_∞
* (q^4;q^5)_∞))"

```

```

    using assms by (simp add: q'_def power_abs power_less_one_iff flip:
qpochhammer_inf_of_real)
    finally show "(( $\lambda j. q \wedge (j^2) / \text{qpochhammer } j \ q \ q$ ) has_sum (1 / ((q;q5)∞
* (q4;q5)∞))) UNIV"
    by (subst (asm) has_sum_of_real_iff)

    have "(( $\lambda j. q' \wedge (j^2 + j) / \text{qpochhammer } j \ q' \ q'$ ) has_sum (1 / ((q'2;q'5)∞
* (q'3;q'5)∞))) UNIV"
    by (rule rogers_ramanujan_complex) auto
    also have "(( $\lambda j. q' \wedge (j^2 + j) / \text{qpochhammer } j \ q' \ q'$ ) = ( $\lambda j. \text{of\_real}$ 
( $q \wedge (j^2+j) / \text{qpochhammer } j \ q \ q$ )))"
    by (simp add: q'_def qpochhammer_of_real)
    also have "(1 / ((q'2;q'5)∞ * (q'3;q'5)∞)) = of_real (1 / ((q2;q5)∞
* (q3;q5)∞)))"
    using assms by (simp add: q'_def power_abs power_less_one_iff flip:
qpochhammer_inf_of_real)
    finally show "(( $\lambda j. q \wedge (j^2+j) / \text{qpochhammer } j \ q \ q$ ) has_sum (1 / ((q2;q5)∞
* (q3;q5)∞))) UNIV"
    by (subst (asm) has_sum_of_real_iff)
qed

unbundle no_qpochhammer_inf_notation

end

```

## References

- [1] G. Andrews and K. Eriksson. *Integer Partitions*. Cambridge University Press, 2004.