

Providing restriction Spaces with an ultrametric Structure

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Abstract

We investigate the relationship between restriction spaces and classical metric structures by instantiating the former as ultrametric spaces. This is classically captured by defining the distance as

$$\text{dist } x \ y = \inf_{x \downarrow n = y \downarrow n} \left(\frac{1}{2} \right)^n$$

but we actually generalize this perspective by introducing a hierarchy of increasingly refined type classes to systematically relate ultrametric and restriction-based notions. This layered approach enables a precise comparison of structural and topological properties. In the end, our main result establishes that completeness in the sense of restriction spaces coincides with standard metric completeness, thus bridging the gap between `Restriction_Spaces` and Banach's fixed-point theorem established in `HOL-Analysis`.

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1 Definitions on Functions of Metric Space

In this theory, we define the notion of lipschitz map, non-expanding map and contraction map. We also establish correspondences.

1.1 Definitions

1.1.1 Lipschitz Map

This notion is a generalization of contraction map and non-expanding map.

definition *lipschitz-with-on* :: $\langle 'a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space}, \text{real}, 'a \text{ set} \rangle \Rightarrow \text{bool}$

where $\langle \text{lipschitz-with-on } f \ \alpha \ A \equiv 0 \leq \alpha \wedge (\forall x \in A. \forall y \in A. \text{dist } (f \ x) \ (f \ y) \leq \alpha * \text{dist } x \ y) \rangle$

abbreviation *lipschitz-with* :: $\langle 'a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space}, \text{real} \rangle \Rightarrow \text{bool}$

where $\langle \text{lipschitz-with } f \ \alpha \equiv \text{lipschitz-with-on } f \ \alpha \ \text{UNIV} \rangle$

lemma *lipschitz-with-onI* :

$\langle \llbracket 0 \leq \alpha; \bigwedge x \ y. \llbracket x \in A; y \in A; x \neq y; f \ x \neq f \ y \rrbracket \implies \text{dist } (f \ x) \ (f \ y) \leq \alpha * \text{dist } x \ y \rrbracket \implies$

$\text{lipschitz-with-on } f \ \alpha \ A \rangle$

unfolding *lipschitz-with-on-def* **by** (*metis dist-eq-0-iff zero-le-dist zero-le-mult-iff*)

lemma *lipschitz-withI* :

$\langle \llbracket 0 \leq \alpha; \bigwedge x \ y. x \neq y \implies f \ x \neq f \ y \implies \text{dist } (f \ x) \ (f \ y) \leq \alpha * \text{dist } x \ y \rrbracket \implies$

$\text{lipschitz-with } f \ \alpha \rangle$

by (rule *lipschitz-with-onI*)

lemma *lipschitz-with-onD1* : $\langle \text{lipschitz-with-on } f \ \alpha \ A \implies 0 \leq \alpha \rangle$
unfolding *lipschitz-with-on-def* **by** *simp*

lemma *lipschitz-withD1* : $\langle \text{lipschitz-with } f \ \alpha \implies 0 \leq \alpha \rangle$
by (rule *lipschitz-with-onD1*)

lemma *lipschitz-with-onD2* :
 $\langle \text{lipschitz-with-on } f \ \alpha \ A \implies x \in A \implies y \in A \implies \text{dist } (f \ x) \ (f \ y) \leq \alpha * \text{dist } x \ y \rangle$
unfolding *lipschitz-with-on-def* **by** *simp*

lemma *lipschitz-withD2* :
 $\langle \text{lipschitz-with } f \ \alpha \implies \text{dist } (f \ x) \ (f \ y) \leq \alpha * \text{dist } x \ y \rangle$
unfolding *lipschitz-with-on-def* **by** *simp*

lemma *lipschitz-with-imp-lipschitz-with-on*: $\langle \text{lipschitz-with } f \ \alpha \implies \text{lipschitz-with-on } f \ \alpha \ A \rangle$
by (*simp add: lipschitz-with-on-def*)

lemma *lipschitz-with-on-imp-lipschitz-with-on-ge* : $\langle \text{lipschitz-with-on } f \ \beta \ A \rangle$
if $\langle \alpha \leq \beta \rangle$ **and** $\langle \text{lipschitz-with-on } f \ \alpha \ A \rangle$
proof (rule *lipschitz-with-onI*)
show $\langle 0 \leq \beta \rangle$ **by** (*meson order-trans lipschitz-with-onD1 that*)
next
fix $x \ y$ **assume** $\langle x \in A \rangle$ **and** $\langle y \in A \rangle$
from $\langle \text{lipschitz-with-on } f \ \alpha \ A \rangle$ [*THEN lipschitz-with-onD2, OF this*]
have $\langle \text{dist } (f \ x) \ (f \ y) \leq \alpha * \text{dist } x \ y \rangle$.
from *order-trans* [*OF this*] **show** $\langle \text{dist } (f \ x) \ (f \ y) \leq \beta * \text{dist } x \ y \rangle$
by (*simp add: mult-right-mono* $\langle \alpha \leq \beta \rangle$)
qed

theorem *lipschitz-with-on-comp-lipschitz-with-on* :
 $\langle \text{lipschitz-with-on } (\lambda x. g \ (f \ x)) \ (\beta * \alpha) \ A \rangle$
if $\langle f \ ' \ A \subseteq B \rangle$ $\langle \text{lipschitz-with-on } g \ \beta \ B \rangle$ $\langle \text{lipschitz-with-on } f \ \alpha \ A \rangle$
proof (rule *lipschitz-with-onI*)
show $\langle 0 \leq \beta * \alpha \rangle$ **by** (*metis lipschitz-with-onD1 mult-nonneg-nonneg that(2, 3)*)
next
fix $x \ y$ **assume** $\langle x \in A \rangle$ **and** $\langle y \in A \rangle$
with $\langle f \ ' \ A \subseteq B \rangle$ **have** $\langle f \ x \in B \rangle$ **and** $\langle f \ y \in B \rangle$ **by** *auto*
have $\langle \text{dist } (g \ (f \ x)) \ (g \ (f \ y)) \leq \beta * \text{dist } (f \ x) \ (f \ y) \rangle$
by (*fact that(2)* [*THEN lipschitz-with-onD2, OF* $\langle f \ x \in B \rangle$ $\langle f \ y \in B \rangle$])
also have $\langle \dots \leq \beta * (\alpha * \text{dist } x \ y) \rangle$

by (fact mult-left-mono[OF that(3)[THEN lipschitz-with-onD2, OF
 $\langle x \in A \rangle \langle y \in A \rangle$]

that(2)[THEN lipschitz-with-onD1]))

also have $\langle \dots = \beta * \alpha * \text{dist } x \ y \rangle$ by simp

finally show $\langle \text{dist } (g (f x)) (g (f y)) \leq \dots \rangle$.

qed

corollary lipschitz-with-comp-lipschitz-with :

$\langle \llbracket \text{lipschitz-with } g \ \beta; \text{lipschitz-with } f \ \alpha \rrbracket \implies$

$\text{lipschitz-with } (\lambda x. g (f x)) (\beta * \alpha) \rangle$

using lipschitz-with-on-comp-lipschitz-with-on by blast

1.1.2 Non-expanding Map

definition non-expanding-on :: $\langle [a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space},$
 $'a \text{ set}] \Rightarrow \text{bool} \rangle$

where $\langle \text{non-expanding-on } f \ A \equiv \text{lipschitz-with-on } f \ 1 \ A \rangle$

abbreviation non-expanding :: $\langle [a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space}]$
 $\Rightarrow \text{bool} \rangle$

where $\langle \text{non-expanding } f \equiv \text{non-expanding-on } f \ \text{UNIV} \rangle$

lemma non-expanding-onI :

$\langle \llbracket \bigwedge x \ y. \llbracket x \in A; y \in A; x \neq y; f x \neq f y \rrbracket \implies \text{dist } (f x) (f y) \leq \text{dist } x \ y \rrbracket \implies$

$\text{non-expanding-on } f \ A \rangle$

by (simp add: lipschitz-with-onI non-expanding-on-def)

lemma non-expandingI :

$\langle \llbracket \bigwedge x \ y. x \neq y \implies f x \neq f y \implies \text{dist } (f x) (f y) \leq \text{dist } x \ y \rrbracket \implies$

$\text{non-expanding } f \rangle$

by (rule non-expanding-onI)

lemma non-expanding-onD :

$\langle \text{non-expanding-on } f \ A \implies x \in A \implies y \in A \implies \text{dist } (f x) (f y) \leq$
 $\text{dist } x \ y \rangle$

by (metis lipschitz-with-onD2 mult-1 non-expanding-on-def)

lemma non-expandingD : $\langle \text{non-expanding } f \implies \text{dist } (f x) (f y) \leq \text{dist } x \ y \rangle$

by (simp add: non-expanding-onD)

lemma non-expanding-imp-non-expanding-on: $\langle \text{non-expanding } f \implies$
 $\text{non-expanding-on } f \ A \rangle$

by (meson non-expandingD non-expanding-onI)

lemma non-expanding-on-comp-non-expanding-on :

$\langle \llbracket f \ 'A \subseteq B; \text{non-expanding-on } g \ B; \text{non-expanding-on } f \ A \rrbracket \implies$

non-expanding-on $(\lambda x. g (f x)) A$
unfolding *non-expanding-on-def*
by (*metis (no-types) lipschitz-with-on-comp-lipschitz-with-on mult-1*)

corollary *non-expanding-comp-non-expanding* :
 $\langle \llbracket \text{non-expanding } g; \text{non-expanding } f \rrbracket \implies \text{non-expanding } (\lambda x. g (f x)) \rangle$
by (*blast intro: non-expanding-on-comp-non-expanding-on*)

1.1.3 Contraction Map

definition *contraction-with-on* :: $\langle ['a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space}, \text{real}, 'a \text{ set}] \Rightarrow \text{bool} \rangle$
where $\langle \text{contraction-with-on } f \alpha A \equiv \alpha < 1 \wedge \text{lipschitz-with-on } f \alpha A \rangle$

abbreviation *contraction-with* :: $\langle ['a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space}, \text{real}] \Rightarrow \text{bool} \rangle$
where $\langle \text{contraction-with } f \alpha \equiv \text{contraction-with-on } f \alpha \text{ UNIV} \rangle$

definition *contraction-on* :: $\langle ['a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space}, 'a \text{ set}] \Rightarrow \text{bool} \rangle$
where $\langle \text{contraction-on } f A \equiv \exists \alpha. \text{contraction-with-on } f \alpha A \rangle$

abbreviation *contraction* :: $\langle ['a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space}] \Rightarrow \text{bool} \rangle$
where $\langle \text{contraction } f \equiv \text{contraction-on } f \text{ UNIV} \rangle$

lemma *contraction-with-onI* :
 $\langle \llbracket 0 \leq \alpha; \alpha < 1; \bigwedge x y. \llbracket x \in A; y \in A; x \neq y; f x \neq f y \rrbracket \implies \text{dist } (f x) (f y) \leq \alpha * \text{dist } x y \rrbracket \implies \text{contraction-with-on } f \alpha A \rangle$
by (*simp add: contraction-with-on-def lipschitz-with-onI*)

lemma *contraction-withI* :
 $\langle \llbracket 0 \leq \alpha; \alpha < 1; \bigwedge x y. x \neq y \implies f x \neq f y \implies \text{dist } (f x) (f y) \leq \alpha * \text{dist } x y \rrbracket \implies \text{contraction-with } f \alpha \rangle$
by (*rule contraction-with-onI*)

lemma *contraction-with-onD1* : $\langle \text{contraction-with-on } f \alpha A \implies 0 \leq \alpha \rangle$
by (*simp add: contraction-with-on-def lipschitz-with-on-def*)

lemma *contraction-withD1* : $\langle \text{contraction-with } f \alpha \implies 0 \leq \alpha \rangle$
by (*simp add: contraction-with-onD1*)

lemma *contraction-with-onD2* : $\langle \text{contraction-with-on } f \alpha A \implies \alpha <$

1)

by (*simp add: contraction-with-on-def lipschitz-with-on-def*)

lemma *contraction-withD2* : $\langle \text{contraction-with } f \ \alpha \implies \alpha < 1 \rangle$
by (*simp add: contraction-with-onD2*)

lemma *contraction-with-onD3* :
 $\langle \text{contraction-with-on } f \ \alpha \ A \implies x \in A \implies y \in A \implies \text{dist } (f \ x) \ (f \ y) \leq \alpha * \text{dist } x \ y \rangle$
by (*simp add: contraction-with-on-def lipschitz-with-on-def*)

lemma *contraction-withD3* : $\langle \text{contraction-with } f \ \alpha \implies \text{dist } (f \ x) \ (f \ y) \leq \alpha * \text{dist } x \ y \rangle$
by (*simp add: contraction-with-on-def lipschitz-withD2*)

lemma *contraction-with-imp-contraction-with-on*:
 $\langle \text{contraction-with } f \ \alpha \implies \text{contraction-with-on } f \ \alpha \ A \rangle$
by (*simp add: contraction-with-on-def lipschitz-with-imp-lipschitz-with-on*)

lemma *contraction-imp-contraction-on*: $\langle \text{contraction } f \implies \text{contraction-on } f \ A \rangle$
using *contraction-on-def contraction-with-imp-contraction-with-on*
by *blast*

lemma *contraction-with-on-imp-contraction-on* :
 $\langle \text{contraction-with-on } f \ \alpha \ A \implies \text{contraction-on } f \ A \rangle$
unfolding *contraction-on-def* **by** *blast*

lemma *contraction-with-imp-contraction*: $\langle \text{contraction-with } f \ \alpha \implies \text{contraction } f \rangle$
by (*simp add: contraction-with-on-imp-contraction-on*)

lemma *contraction-onE*:
 $\langle \llbracket \text{contraction-on } f \ A; \bigwedge \alpha. \text{contraction-with-on } f \ \alpha \ A \implies \text{thesis} \rrbracket \implies \text{thesis} \rangle$
unfolding *contraction-on-def* **by** *blast*

lemma *contractionE*:
 $\langle \llbracket \text{contraction } f; \bigwedge \alpha. \text{contraction-with } f \ \alpha \implies \text{thesis} \rrbracket \implies \text{thesis} \rangle$
by (*elim contraction-onE*)

lemma *contraction-with-on-imp-contraction-with-on-ge* :
 $\langle \llbracket \alpha \leq \beta; \beta < 1; \text{contraction-with-on } f \ \alpha \ A \rrbracket \implies \text{contraction-with-on } f \ \beta \ A \rangle$
by (*simp add: contraction-with-on-def lipschitz-with-on-imp-lipschitz-with-on-ge*)

1.2 Properties

lemma *contraction-with-on-imp-lipschitz-with-on*[simp] :
 $\langle \text{contraction-with-on } f \ \alpha \ A \implies \text{lipschitz-with-on } f \ \alpha \ A \rangle$
by (simp add: contraction-with-on-def)

lemma *non-expanding-on-imp-lipschitz-with-one-on*[simp] :
 $\langle \text{non-expanding-on } f \ A \implies \text{lipschitz-with-on } f \ 1 \ A \rangle$
by (simp add: non-expanding-on-def)

lemma *contraction-on-imp-non-expanding-on*[simp] :
 $\langle \text{contraction-on } f \ A \implies \text{non-expanding-on } f \ A \rangle$

proof (elim contraction-onE, rule non-expanding-onI)

fix $\alpha \ x \ y$ **assume** $\langle x \in A \rangle$ **and** $\langle y \in A \rangle$ **and** contra : $\langle \text{contraction-with-on } f \ \alpha \ A \rangle$

show $\langle \text{dist } (f \ x) \ (f \ y) \leq \text{dist } x \ y \rangle$

by (rule order-trans[OF contra[THEN contraction-with-onD3], OF $\langle x \in A \rangle \ \langle y \in A \rangle$]])

(metis contra contraction-with-on-def mult-less-cancel-right1 nle-le order-less-le zero-le-dist)

qed

lemma *contraction-with-on-comp-contraction-with-on* :

$\langle \text{contraction-with-on } (\lambda x. g \ (f \ x)) \ (\beta * \alpha) \ A \rangle$

if $\langle f \ ' \ A \subseteq B \rangle$ $\langle \text{contraction-with-on } g \ \beta \ B \rangle$ $\langle \text{contraction-with-on } f \ \alpha \ A \rangle$

proof (unfold contraction-with-on-def, intro conjI)

from that(2, 3)[THEN contraction-with-onD2] that(2)[THEN contraction-with-onD1]

show $\langle \beta * \alpha < 1 \rangle$ **by** (metis dual-order.strict-trans2 mult-less-cancel-left1 nle-le order-less-le)

next

show $\langle \text{lipschitz-with-on } (\lambda x. g \ (f \ x)) \ (\beta * \alpha) \ A \rangle$

by (rule lipschitz-with-on-comp-lipschitz-with-on[OF $\langle f \ ' \ A \subseteq B \rangle$])
(simp-all add: that(2, 3))

qed

corollary *contraction-with-comp-contraction-with* :

$\langle \llbracket \text{contraction-with } g \ \beta; \text{contraction-with } f \ \alpha \rrbracket \implies \text{contraction-with } (\lambda x. g \ (f \ x)) \ (\beta * \alpha) \rangle$

by (blast intro: contraction-with-on-comp-contraction-with-on)

corollary *contraction-on-comp-contraction-on* :

$\langle \llbracket f \ ' \ A \subseteq B; \text{contraction-on } g \ B; \text{contraction-on } f \ A \rrbracket \implies \text{contraction-on } (\lambda x. g \ (f \ x)) \ A \rangle$

proof (elim contraction-onE)

fix $\alpha \ \beta$ **assume** $\langle f \ ' \ A \subseteq B \rangle$ $\langle \text{contraction-with-on } g \ \beta \ B \rangle$ $\langle \text{contraction-with-on } f \ \alpha \ A \rangle$

from contraction-with-on-comp-contraction-with-on[OF this]

show $\langle \text{contraction-on } (\lambda x. g (f x)) A \rangle$ **by** (fact contraction-with-on-imp-contraction-on)
qed

corollary *contraction-comp-contraction* :

$\langle \llbracket \text{contraction } g; \text{contraction } f \rrbracket \implies \text{contraction } (\lambda x. g (f x)) \rangle$
by (blast intro: contraction-on-comp-contraction-on)

lemma *contraction-with-on-comp-non-expanding-on* :

$\langle \text{contraction-with-on } (\lambda x. g (f x)) \beta A \rangle$
if $\langle f ' A \subseteq B \rangle \langle \text{contraction-with-on } g \beta B \rangle \langle \text{non-expanding-on } f A \rangle$
proof (unfold contraction-with-on-def, intro conjI)
from that(2)[THEN contraction-with-onD2] **show** $\langle \beta < 1 \rangle$.
next
show $\langle \text{lipschitz-with-on } (\lambda x. g (f x)) \beta A \rangle$
by (rule lipschitz-with-on-comp-lipschitz-with-on[OF $\langle f ' A \subseteq B \rangle$,
of $g \beta 1$, simplified])
(simp-all add: that(2, 3))
qed

corollary *contraction-with-comp-non-expanding* :

$\langle \llbracket \text{contraction-with } g \beta; \text{non-expanding } f \rrbracket \implies \text{contraction-with } (\lambda x. g (f x)) \beta \rangle$
by (blast intro: contraction-with-on-comp-non-expanding-on)

corollary *contraction-on-comp-non-expanding-on* :

$\langle \llbracket f ' A \subseteq B; \text{contraction-on } g B; \text{non-expanding-on } f A \rrbracket \implies \text{contraction-on } (\lambda x. g (f x)) A \rangle$
by (metis contraction-on-def contraction-with-on-comp-non-expanding-on)

corollary *contraction-comp-non-expanding* :

$\langle \llbracket \text{contraction } g; \text{non-expanding } f \rrbracket \implies \text{contraction } (\lambda x. g (f x)) \rangle$
by (blast intro: contraction-on-comp-non-expanding-on)

lemma *non-expanding-on-comp-contraction-with-on* :

$\langle \text{contraction-with-on } (\lambda x. g (f x)) \alpha A \rangle$
if $\langle f ' A \subseteq B \rangle \langle \text{non-expanding-on } g B \rangle \langle \text{contraction-with-on } f \alpha A \rangle$
proof (unfold contraction-with-on-def, intro conjI)
from that(3)[THEN contraction-with-onD2] **show** $\langle \alpha < 1 \rangle$.
next
show $\langle \text{lipschitz-with-on } (\lambda x. g (f x)) \alpha A \rangle$
by (rule lipschitz-with-on-comp-lipschitz-with-on[OF $\langle f ' A \subseteq B \rangle$,
of $g 1$, simplified])
(simp-all add: that(2, 3))
qed

corollary *non-expanding-comp-contraction-with* :

$\langle \llbracket \text{non-expanding } g; \text{contraction-with } f \alpha \rrbracket \implies \text{contraction-with } (\lambda x.$

$g (f x) \alpha$
by (*blast intro: non-expanding-on-comp-contraction-with-on*)

corollary *non-expanding-on-comp-contraction-on* :
 $\langle \llbracket f ' A \subseteq B; \text{non-expanding-on } g B; \text{contraction-on } f A \rrbracket \implies \text{contraction-on } (\lambda x. g (f x)) A \rangle$
by (*metis contraction-on-def non-expanding-on-comp-contraction-with-on*)

corollary *non-expanding-comp-contraction* :
 $\langle \llbracket \text{non-expanding } g; \text{contraction } f \rrbracket \implies \text{contraction } (\lambda x. g (f x)) \rangle$
by (*blast intro: non-expanding-on-comp-contraction-on*)

1.3 Banach's fixed-point Theorems

We rewrite the Banach's fixed-point theorems with our new definition.

theorem *Banach-fix-type* : $\langle \text{contraction } f \implies \exists !x. f x = x \rangle$
for $f :: \langle 'a :: \text{complete-space} \Rightarrow 'a \rangle$
by (*elim contractionE*)
(metis banach-fix-type contraction-withD1 contraction-withD2 contraction-withD3)

theorem *Banach-fix*:
 $\langle \text{contraction-on } f s \implies \exists !x. x \in s \wedge f x = x \rangle$ **if** $\langle \text{complete } s \rangle \langle s \neq \{\} \rangle \langle f ' s \subseteq s \rangle$

proof (*elim contraction-onE, intro Banach-fix[OF $\langle \text{complete } s \rangle \langle s \neq \{\} \rangle - - \langle f ' s \subseteq s \rangle$] ballI*)

show $\langle \text{contraction-with-on } f \alpha s \implies 0 \leq \alpha \rangle$ **for** α **by** (*fact contraction-with-onD1*)

next

show $\langle \text{contraction-with-on } f \alpha s \implies \alpha < 1 \rangle$ **for** α **by** (*fact contraction-with-onD2*)

next

show $\langle \text{contraction-with-on } f \alpha s \implies x \in s \implies y \in s \implies \text{dist } (f x) (f y) \leq \alpha * \text{dist } x y \rangle$ **for** $\alpha x y$ **by** (*fact contraction-with-onD3*)

qed

2 Locales factorizing the proof Work

2.1 Preliminaries on strictly decreasing Sequences

abbreviation *strict-decseq* :: $\langle (\text{nat} \Rightarrow 'a :: \text{order}) \Rightarrow \text{bool} \rangle$

where $\langle \text{strict-decseq} \equiv \text{monotone } (<) (\lambda x y. y < x) \rangle$

lemma *strict-decseq-def* : $\langle \text{strict-decseq } X \longleftrightarrow (\forall m n. m < n \longrightarrow X) \rangle$

$n < X m$ ›
by (*fact monotone-def*)

lemma *strict-decseqI* : ‹*strict-decseq X*› **if** ‹ $\bigwedge n. X (Suc\ n) < X\ n$ ›
by (*metis Suc-le-eq decseqD decseq-Suc-iff le-less-trans nless-le strict-decseq-def that*)

lemma *strict-decseqD* : ‹*strict-decseq X* $\implies m < n \implies X\ n < X\ m$ ›
using *strict-decseq-def* **by** *blast*

lemma *strict-decseq-def-bis* : ‹*strict-decseq X* $\longleftrightarrow (\forall m\ n. X\ n < X\ m \longleftrightarrow m < n)$ ›
by (*metis linorder-less-linear order-less-imp-not-less strict-decseq-def*)

lemma *strict-decseq-def-ter* : ‹*strict-decseq X* $\longleftrightarrow (\forall m\ n. X\ n \leq X\ m \longleftrightarrow m \leq n)$ ›
unfolding *strict-decseq-def-bis*
by (*rule iffI, metis le-simps(2) order-le-less, simp add: less-le-not-le*)

lemma *strict-decseq-imp-decseq* : ‹*strict-decseq* $\sigma \implies decseq\ \sigma$ ›
by (*simp add: monotone-on-def order-le-less*)

lemma *strict-decseq-SucI* : ‹ $(\bigwedge n. X (Suc\ n) < X\ n) \implies strict-decseq\ X$ ›
by (*metis Suc-le-eq decseqD decseq-Suc-iff less-le-not-le order-less-le strict-decseq-def*)

lemma *strict-decseq-SucD* : ‹*strict-decseq A* $\implies A (Suc\ i) < A\ i$ ›
by (*simp add: strict-decseq-def*)

Classically, a restriction space is given the structure of a metric space by defining $dist\ x\ y \equiv Inf\ \{(1 / 2) \wedge n \mid n. x \downarrow n = y \downarrow n\}$. This obviously also works if we replace $1 / 2$ by any real δ such that $0 < \delta$ and $\delta < 1$. But more generally, this still works if we set $dist\ x\ y \equiv Inf\ \{\sigma\ n \mid n. x \downarrow n = y \downarrow n\}$ where σ is a sequence of *real* verifying $\forall n. 0 < \sigma\ n$ and $\sigma \longrightarrow 0$. As you would expect, the more structure you have, the more powerful theorems you get. We explore all these variants in the theory below.

2.2 The Construction with Locales

Our formalization will extend the class *metric-space*. But some proofs are redundant, especially when it comes to the product type. So first we will be working with locales, and interpret them with the classes.

locale *NonDecseqRestrictionSpace* = *PreorderRestrictionSpace* +
 — Factorization of the proof work.
fixes $M :: \langle 'a \text{ set} \rangle$ **and** $\text{restriction-}\sigma :: \langle \text{nat} \Rightarrow \text{real} \rangle (\sigma_{\downarrow})$
and $\text{restriction-dist} :: \langle 'a \Rightarrow 'a \Rightarrow \text{real} \rangle (\text{dist}_{\downarrow})$
assumes $\text{restriction-}\sigma\text{-tendsto-zero} : \langle \text{restriction-}\sigma \longrightarrow 0 \rangle$
and $\text{zero-less-restriction-}\sigma \text{ [simp]} : \langle 0 < \text{restriction-}\sigma \ n \rangle$
and $\text{dist-restriction-is} :$
 $\langle \text{dist}_{\downarrow} \ x \ y = (\text{INF } n \in \text{restriction-related-set } x \ y. \text{restriction-}\sigma \ n) \rangle$
begin

lemma $\text{zero-le-restriction-}\sigma \text{ [simp]} : \langle 0 \leq \sigma_{\downarrow} \ n \rangle$
by (*simp add: order-less-imp-le*)

lemma $\text{restriction-}\sigma\text{-neq-zero [simp]} : \langle \sigma_{\downarrow} \ n \neq 0 \rangle$
by (*metis zero-less-restriction-}\sigma \text{ order-less-irrefl}*)

lemma $\text{bounded-range-restriction-}\sigma : \langle \text{bounded } (\text{range } \sigma_{\downarrow}) \rangle$
by (*fact convergent-imp-bounded[OF restriction-}\sigma\text{-tendsto-zero]*)

abbreviation $\text{restriction-}\sigma\text{-related-set} :: \langle 'a \Rightarrow 'a \Rightarrow \text{real set} \rangle$
where $\langle \text{restriction-}\sigma\text{-related-set } x \ y \equiv \sigma_{\downarrow} \text{ `restriction-related-set } x \ y \rangle$

abbreviation $\text{restriction-}\sigma\text{-not-related-set} :: \langle 'a \Rightarrow 'a \Rightarrow \text{real set} \rangle$
where $\langle \text{restriction-}\sigma\text{-not-related-set } x \ y \equiv \sigma_{\downarrow} \text{ `restriction-not-related-set } x \ y \rangle$

lemma $\text{nonempty-restriction-}\sigma\text{-related-set} :$
 $\langle \text{restriction-}\sigma\text{-related-set } x \ y \neq \{\} \rangle$ **by** *simp*

lemma $\text{restriction-}\sigma\text{-related-set-}\cup\text{-restriction-}\sigma\text{-not-related-set} :$
 $\langle \text{restriction-}\sigma\text{-related-set } x \ y \cup \text{restriction-}\sigma\text{-not-related-set } x \ y = \text{range } \sigma_{\downarrow} \rangle$
by *blast*

lemma $\langle \text{bdd-above } (\text{restriction-}\sigma\text{-related-set } x \ y) \rangle$
by (*meson bdd-above.I2 bdd-above.unfold bounded-imp-bdd-above bounded-range-restriction-}\sigma \text{ rangeI}*)

lemma $\langle \text{bdd-above } (\text{restriction-}\sigma\text{-not-related-set } x \ y) \rangle$
by (*meson bdd-above.E bdd-above.I2 bounded-imp-bdd-above bounded-range-restriction-}\sigma \text{ range-eqI}*)

lemma $\text{bounded-restriction-}\sigma\text{-related-set} : \langle \text{bounded } (\text{restriction-}\sigma\text{-related-set } x \ y) \rangle$

by (*meson bounded-range-restriction- σ bounded-subset image-mono top-greatest*)

lemma *bounded-restriction- σ -not-related-set*: $\langle \text{bounded } (\text{restriction-}\sigma\text{-not-related-set } x \ y) \rangle$

by (*meson bounded-range-restriction- σ bounded-subset image-mono subset-UNIV*)

corollary *restriction-space-Inf-properties*:

$\langle a \in \text{restriction-}\sigma\text{-related-set } x \ y \implies \text{dist}_\downarrow x \ y \leq a \rangle$

$\langle \llbracket \bigwedge a. a \in \text{restriction-}\sigma\text{-related-set } x \ y \implies b \leq a \rrbracket \implies b \leq \text{dist}_\downarrow x \ y \rangle$

unfolding *dist-restriction-is*

by (*simp-all add: bounded-has-Inf(1) bounded-restriction- σ -related-set cInf-greatest*)

lemma *restriction- σ -related-set-alt* :

$\langle \text{restriction-}\sigma\text{-related-set } x \ y = \{ \sigma_\downarrow n \mid n. n \in \text{restriction-related-set } x \ y \} \rangle$

by *blast*

lemma *exists-less-restriction- σ* : $\langle \exists n. m < n \wedge \sigma_\downarrow n < \sigma_\downarrow m \rangle$

proof (*rule ccontr*)

assume $\langle \neg (\exists n > m. \sigma_\downarrow n < \sigma_\downarrow m) \rangle$

hence $\langle \forall n \geq m. \sigma_\downarrow m \leq \sigma_\downarrow n \rangle$

by (*metis linorder-not-le nle-le*)

hence $\langle \neg \sigma_\downarrow \longrightarrow 0 \rangle$

by (*meson Lim-bounded2 linorder-not-le zero-less-restriction- σ*)

with *restriction- σ -tendsto-zero* **show** *False* **by** *simp*

qed

lemma $\langle \text{dist}_\downarrow x \ y = \text{Inf } (\text{restriction-}\sigma\text{-related-set } x \ y) \rangle$

by (*fact dist-restriction-is*)

lemma *not-related-imp-dist-restriction-is-some-restriction- σ* :

$\langle \exists n. \text{dist}_\downarrow x \ y = \sigma_\downarrow n \wedge (\forall m \leq n. x \downarrow m \not\approx y \downarrow m) \rangle$ **if** $\langle \neg x \lesssim y \rangle$

proof –

have $\langle \text{finite } (\text{restriction-related-set } x \ y) \rangle$

by (*simp add: finite-restriction-related-set-iff* $\langle \neg x \lesssim y \rangle$)

have $\langle \text{Inf } (\text{restriction-}\sigma\text{-related-set } x \ y) \in \text{restriction-}\sigma\text{-related-set } x \ y \rangle$

by (*rule closed-contains-Inf[OF nonempty-restriction- σ -related-set]*)

(*simp-all add: finite (restriction-related-set x y) finite-imp-closed*)

hence $\langle \text{dist}_\downarrow x \ y \in \text{restriction-}\sigma\text{-related-set } x \ y \rangle$

by (*fold dist-restriction-is*)

with *restriction-related-le* **obtain** *n*

where $\langle n \in \text{restriction-related-set } x \ y \ \langle \text{dist}_\downarrow x \ y = \sigma_\downarrow n \rangle \rangle$

$\langle \forall m \leq n. x \downarrow m \lesssim y \downarrow m \rangle$ **by** *blast*
with $\langle \text{dist}_{\downarrow} x y \in \text{restriction-}\sigma\text{-related-set } x y \rangle$ **show** *?thesis* **by** *blast*
qed

lemma *not-related-imp-dist-restriction-le-some-restriction- σ* :
 $\langle \neg x \lesssim y \implies \exists n. \text{dist}_{\downarrow} x y \leq \sigma_{\downarrow} n \wedge (\neg x \downarrow \text{Suc } n \lesssim y \downarrow \text{Suc } n) \wedge (\forall m \leq n. x \downarrow m \lesssim y \downarrow m) \rangle$
by (*blast intro: restriction-space-Inf-properties*
dest: ex-not-restriction-related-optimized)

lemma *restriction-dist-eq-0-iff-related* : $\langle \text{dist}_{\downarrow} x y = 0 \longleftrightarrow x \lesssim y \rangle$
proof (*rule iffI*)
show $\langle \text{dist}_{\downarrow} x y = 0 \implies x \lesssim y \rangle$
by (*erule contrapos-pp*)
(auto dest: not-related-imp-dist-restriction-is-some-restriction- σ)
next
show $\langle \text{dist}_{\downarrow} x y = 0 \rangle$ **if** $\langle x \lesssim y \rangle$
proof (*rule order-antisym*)
show $\langle 0 \leq \text{dist}_{\downarrow} x y \rangle$ **by** (*simp add: cINF-greatest dist-restriction-is*)
next
define Δ **where** $\langle \Delta n \equiv - \sigma_{\downarrow} n \rangle$ **for** n
have $*$: $\langle \Delta n \leq - \text{dist}_{\downarrow} x y \rangle$ **for** n
unfolding Δ -*def* **using** $\langle x \lesssim y \rangle$ *restriction-space-Inf-properties(1)*
by *simp (metis UNIV-restriction-related-set-iff rangeI)*
from *restriction- σ -tendsto-zero tendsto-minus-cancel-left*
have $\langle \Delta \longrightarrow 0 \rangle$ **unfolding** Δ -*def* **by** *force*
from *lim-le[OF convergentI[OF $\langle \Delta \longrightarrow 0 \rangle$] *] limI[OF $\langle \Delta \longrightarrow 0 \rangle$]*
show $\langle \text{dist}_{\downarrow} x y \leq 0 \rangle$ **by** *simp*
qed
qed

end

locale *DecseqRestrictionSpace* = *NonDecseqRestrictionSpace* +
assumes *decseq-restriction- σ* : $\langle \text{decseq } \sigma_{\downarrow} \rangle$
begin

lemma *dist-restriction-is-bis* :
 $\langle \text{dist}_{\downarrow} x y = (\text{if } x \lesssim y \text{ then } 0 \text{ else } \sigma_{\downarrow} (\text{Sup } (\text{restriction-related-set } x y))) \rangle$
if $\langle x \in M \rangle$ **and** $\langle y \in M \rangle$
proof (*split if-split, intro conjI impI*)
from $\langle x \in M \rangle$ **and** $\langle y \in M \rangle$ **show** $\langle x \lesssim y \implies \text{restriction-dist } x y = 0 \rangle$

```

    by (simp add: restriction-dist-eq-0-iff-related)
  next
  show  $\langle \text{dist}_\downarrow x y = \sigma_\downarrow (\text{Sup} (\text{restriction-related-set } x y)) \rangle$  if  $\langle \neg x \lesssim y \rangle$ 
  proof (rule order-antisym)
    show  $\langle \text{dist}_\downarrow x y \leq \sigma_\downarrow (\text{Sup} (\text{restriction-related-set } x y)) \rangle$ 
    unfolding dist-restriction-is
    by (rule cINF-lower[OF - Sup-in-restriction-related-set[OF  $\langle \neg x \lesssim y \rangle$ ]])
    (meson bdd-below.I2 zero-le-restriction- $\sigma$ )
  next
  show  $\langle \sigma_\downarrow (\text{Sup} (\text{restriction-related-set } x y)) \leq \text{dist}_\downarrow x y \rangle$ 
  proof (unfold dist-restriction-is, rule cINF-greatest)
  show  $\langle \text{restriction-related-set } x y \neq \{\} \rangle$  by (fact nonempty-restriction-related-set)
  next
  fix n assume  $\langle n \in \text{restriction-related-set } x y \rangle$ 
  hence  $\langle n \leq \text{Sup} (\text{restriction-related-set } x y) \rangle$ 
  by (metis  $\langle n \in \text{restriction-related-set } x y \rangle$  le-cSup-finite
    finite-restriction-related-set-iff  $\langle \neg x \lesssim y \rangle$ )
  from decseqD[OF decseq-restriction- $\sigma$  this]
  show  $\langle \text{restriction-}\sigma (\text{Sup} (\text{restriction-related-set } x y)) \leq \text{restriction-}\sigma n \rangle$  .
  qed
  qed
  qed

```

lemma *not-eq-imp-dist-restriction-is-restriction- σ -Sup-restriction-eq* :

```

 $\langle \text{dist}_\downarrow x y = \sigma_\downarrow (\text{Sup} (\text{restriction-related-set } x y)) \rangle$ 
 $\langle \forall m \leq \text{Sup} (\text{restriction-related-set } x y). x \downarrow m \lesssim y \downarrow m \rangle$ 
 $\langle \forall m > \text{Sup} (\text{restriction-related-set } x y). \neg x \downarrow m \lesssim y \downarrow m \rangle$ 
if  $\langle \neg x \lesssim y \rangle$  and  $\langle x \in M \rangle$  and  $\langle y \in M \rangle$ 
subgoal by (subst dist-restriction-is-bis; simp add:  $\langle \neg x \lesssim y \rangle \langle x \in M \rangle \langle y \in M \rangle$ )
subgoal using Sup-in-restriction-related-set[OF  $\langle \neg x \lesssim y \rangle$ ] restriction-related-le by blast
using cSup-upper[OF - bdd-above-restriction-related-set-iff[THEN iffD2, OF  $\langle \neg x \lesssim y \rangle$ ],
of  $\langle \text{Suc} (\text{Sup} (\text{restriction-related-set } x y)) \rangle$ ]
by (metis (mono-tags, lifting) Suc-leI dual-order.refl mem-Collect-eq not-less-eq-eq restriction-related-le)

```

theorem *restriction-dist-tendsto-zero-independent-of-restriction- σ* :

- Very powerful theorem: the convergence of the distance to 0 is actually independent from the restriction sequence chosen.
- This is the theorem for which we had to work with locales first.

assumes $\langle \text{DecseqRestrictionSpace } (\downarrow) (\lesssim) \text{ restriction-}\sigma' \text{ restriction-dist}' \rangle$
and $\langle \Sigma \in M \rangle$ **and** $\langle \text{range } \sigma \subseteq M \rangle$
shows $\langle (\lambda n. \text{dist}_{\downarrow} (\sigma n) \Sigma) \longrightarrow 0 \iff (\lambda n. \text{restriction-dist}' (\sigma n) \Sigma) \longrightarrow 0 \rangle$
proof –
 { **fix** $\text{restriction-}\sigma \text{ restriction-dist } \text{restriction-}\sigma' \text{ restriction-dist}'$
 assume $a1 : \langle \text{DecseqRestrictionSpace } (\downarrow) (\lesssim) \text{ restriction-}\sigma \text{ restriction-dist} \rangle$
 and $a2 : \langle \text{DecseqRestrictionSpace } (\downarrow) (\lesssim) \text{ restriction-}\sigma' \text{ restriction-dist}' \rangle$
 and $*$: $\langle (\lambda n. \text{restriction-dist} (\sigma n) \Sigma) \longrightarrow 0 \rangle$

 interpret $i1 : \text{DecseqRestrictionSpace } \langle (\downarrow) \rangle \langle (\lesssim) \rangle M \text{ restriction-}\sigma \text{ restriction-dist}$ **by** (fact a1)
 interpret $i2 : \text{DecseqRestrictionSpace } \langle (\downarrow) \rangle \langle (\lesssim) \rangle M \text{ restriction-}\sigma' \text{ restriction-dist}'$ **by** (fact a2)

 have $\langle (\lambda n. \text{restriction-dist}' (\sigma n) \Sigma) \longrightarrow 0 \rangle$
 proof (rule metric-LIMSEQ-I)
 fix $\varepsilon :: \text{real}$ **assume** $\langle 0 < \varepsilon \rangle$

 from $\text{metric-LIMSEQ-D}[OF i2.\text{restriction-}\sigma\text{-tendsto-zero } \langle 0 < \varepsilon \rangle]$
 obtain N **where** $** : \langle N \leq n \implies \text{restriction-}\sigma' n < \varepsilon \rangle$ **for** n
 by *auto*

 fix $N' :: \text{nat}$

 have $\langle \exists N'. \forall n \geq N'. N \in \text{restriction-related-set } (\sigma n) \Sigma \rangle$
 proof (rule ccontr)
 assume $\langle \nexists N'. \forall n \geq N'. N \in \text{restriction-related-set } (\sigma n) \Sigma \rangle$
 hence $*** : \langle \forall N'. \exists n \geq N'. \neg \sigma n \downarrow N \lesssim \Sigma \downarrow N \rangle$ **by** *simp*
 have $**** : \langle \forall N'. \exists n \geq N'. \text{restriction-}\sigma N \leq \text{restriction-dist} (\sigma n) \Sigma \rangle$
 proof (rule allI)
 fix $N' :: \text{nat}$
 from $***$ **obtain** n **where** $\langle N' \leq n \rangle \langle \neg \sigma n \downarrow N \lesssim \Sigma \downarrow N \rangle$
 by *blast*
 hence $\langle \neg \sigma n \lesssim \Sigma \rangle$ **using** *mono-restriction-related* **by** *blast*
 have $\langle \text{restriction-}\sigma N \leq \text{restriction-dist} (\sigma n) \Sigma \rangle$
 proof (*subst i1.dist-restriction-is-bis*)
 show $\langle \sigma n \in M \rangle \langle \Sigma \in M \rangle$
 by (*simp-all add: assms(2, 3) range-subsetD*)
 next
 show $\langle \text{restriction-}\sigma N \leq (\text{if } \sigma n \lesssim \Sigma \text{ then } 0 \text{ else } \text{restriction-}\sigma (\text{Sup } (\text{restriction-related-set } (\sigma n) \Sigma))) \rangle$
 using $\langle \neg \sigma n \lesssim \Sigma \rangle \langle \neg \sigma n \downarrow N \lesssim \Sigma \downarrow N \rangle$ *assms(2, 3)*

```

nle-le
  not-eq-imp-dist-restriction-is-restriction- $\sigma$ -Sup-restriction-eq(2)
  by (fastforce intro: decseqD[OF i1.decseq-restriction- $\sigma$ ])
  qed
  with  $\langle N' \leq n \rangle$  show  $\langle \exists n \geq N'. \text{restriction-}\sigma N \leq \text{restriction-dist}$ 
  ( $\sigma n$ )  $\Sigma \rangle$  by blast
  qed
  from metric-LIMSEQ-D[OF *, of  $\langle \text{restriction-}\sigma N \rangle$ ]
  have  $\langle \exists N''. \forall n \geq N''. \text{restriction-dist} (\sigma n) \Sigma < \text{restriction-}\sigma$ 
   $N \rangle$ 
  by (metis abs-of-nonneg i1.zero-le-restriction- $\sigma$  i1.zero-less-restriction- $\sigma$ 
  norm-conv-dist order-trans real-norm-def
  verit-comp-simplify1(3))
  with **** show False by fastforce
  qed

  then obtain  $N'$  where **** :  $\langle N' \leq n \implies N \in \text{restriction-}$ 
   $\text{related-set} (\sigma n) \Sigma \rangle$  for  $n$  by blast
  have  $\langle \text{restriction-dist}' (\sigma n) \Sigma < \varepsilon \rangle$  if  $\langle N' \leq n \rangle$  for  $n$ 
  proof (rule le-less-trans[OF i2.restriction-space-Inf-properties(1),
  of  $\langle \text{restriction-}\sigma' N \rangle$ ])
    from  $\langle N' \leq n \rangle$  ****
    show  $\langle \text{restriction-}\sigma' N \in i2.\text{restriction-}\sigma\text{-related-set} (\sigma n) \Sigma \rangle$ 
  by blast
  next
    show  $\langle \text{restriction-}\sigma' N < \varepsilon \rangle$  by (simp add: **)
  qed
  thus  $\langle \exists N. \forall n \geq N. \text{dist} (\text{restriction-dist}' (\sigma n) \Sigma) 0 < \varepsilon \rangle$ 
  by (metis abs-of-nonneg dist-0-norm dist-commute i2.not-related-imp-dist-restriction-is-some-restric
  i2.restriction-dist-eq-0-iff-related i2.zero-less-restriction- $\sigma$ 
  order-less-imp-not-less real-norm-def
  verit-comp-simplify1(3))
  qed }
  note * = this

  show  $\langle (\lambda n. \text{dist}_\downarrow (\sigma n) \Sigma) \longrightarrow 0 = (\lambda n. \text{restriction-dist}' (\sigma n)$ 
   $\Sigma) \longrightarrow 0 \rangle$ 
  using * DecseqRestrictionSpace-axioms assms(1) by blast
  qed

end

```

3 Ultrametric Structure of restriction Spaces

This has only be proven with the sort constraint, not inside the context of the class *metric-space* ...

context *metric-space* **begin**

lemma *LIMSEQ-def* : $\langle X \longrightarrow L \longleftrightarrow (\forall r > 0. \exists no. \forall n \geq no. \text{dist } (X \ n) \ L < r) \rangle$

unfolding *tendsto-iff eventually-sequentially ..*

lemma *LIMSEQ-iff-nz*: $\langle X \longrightarrow L \longleftrightarrow (\forall r > 0. \exists no > 0. \forall n \geq no. \text{dist } (X \ n) \ L < r) \rangle$

by (*meson Suc-leD LIMSEQ-def zero-less-Suc*)

lemma *metric-LIMSEQ-I*: $\langle (\bigwedge r. 0 < r \implies \exists no. \forall n \geq no. \text{dist } (X \ n) \ L < r) \implies X \longrightarrow L \rangle$

by (*simp add: LIMSEQ-def*)

lemma *metric-LIMSEQ-D*: $\langle X \longrightarrow L \implies 0 < r \implies \exists no. \forall n \geq no. \text{dist } (X \ n) \ L < r \rangle$

by (*simp add: LIMSEQ-def*)

lemma *LIMSEQ-dist-iff*:

$\langle f \longrightarrow l \longleftrightarrow (\lambda x. \text{dist } (f \ x) \ l \longrightarrow 0) \rangle$

proof (*unfold LIMSEQ-def, rule iffI*)

show $\langle (\lambda n. \text{dist } (f \ n) \ l \longrightarrow 0 \implies (\forall r > 0. \exists no. \forall n \geq no. \text{dist } (f \ n) \ l < r) \rangle$

by (*metis (mono-tags, lifting) eventually-at-top-linorder order-tendstoD(2)*)

next

show $\langle \forall r > 0. \exists no. \forall n \geq no. \text{dist } (f \ n) \ l < r \implies (\lambda n. \text{dist } (f \ n) \ l \longrightarrow 0) \rangle$

by (*simp add: metric-space-class.LIMSEQ-def*)

qed

lemma *Cauchy-converges-subseq*:

fixes $u :: \text{nat} \Rightarrow 'a$

assumes *Cauchy u*

strict-mono r

$(u \circ r) \longrightarrow l$

shows $u \longrightarrow l$

proof –

have *: *eventually* $(\lambda n. \text{dist } (u \ n) \ l < e)$ *sequentially* **if** $e > 0$ **for** e

proof –

have $e/2 > 0$ **using** *that* **by** *auto*

then obtain $N1$ **where** $N1: \bigwedge m \ n. m \geq N1 \implies n \geq N1 \implies \text{dist } (u \ m) \ (u \ n) < e/2$

using $\langle \text{Cauchy } u \rangle$ **unfolding** *Cauchy-def* **by** *blast*

obtain $N2$ **where** $N2: \bigwedge n. n \geq N2 \implies \text{dist } ((u \circ r) \ n) \ l < e / 2$

using *order-tendstoD(2)[OF iffD1[OF LIMSEQ-dist-iff] $\langle (u \circ r) \longrightarrow l \rangle$ $\langle e/2 > 0 \rangle$]*

```

    unfolding eventually-sequentially by auto
have  $\text{dist } (u \ n) \ l < e$  if  $n \geq \max \ N1 \ N2$  for  $n$ 
proof –
    have  $\text{dist } (u \ n) \ l \leq \text{dist } (u \ n) \ ((u \circ r) \ n) + \text{dist } ((u \circ r) \ n) \ l$ 
      by (rule dist-triangle)
    also have  $\dots < e/2 + e/2$ 
    proof (intro add-strict-mono)
      show  $\text{dist } (u \ n) \ ((u \circ r) \ n) < e / 2$ 
        using  $N1[\text{of } n \ r \ n] \ N2[\text{of } n]$  that unfolding comp-def
      by (meson assms(2) le-trans max.bounded-iff strict-mono-imp-increasing)
      show  $\text{dist } ((u \circ r) \ n) \ l < e / 2$ 
        using  $N2$  that by auto
      qed
    finally show ?thesis by simp
  qed
then show ?thesis unfolding eventually-sequentially by blast
qed
have  $(\lambda n. \text{dist } (u \ n) \ l) \longrightarrow 0$ 
  by (simp add: less-le-trans * order-tendstoI)
then show ?thesis using LIMSEQ-dist-iff by auto
qed

```

end

3.1 The Construction with Classes

```

class restriction- $\sigma$  = restriction-space +
  fixes restriction- $\sigma$  ::  $\langle 'a \ \textit{itself} \Rightarrow \text{nat} \Rightarrow \text{real} \rangle \ (\langle \sigma \downarrow \rangle)$ 

```

```

setup  $\langle \textit{Sign.add-const-constraint} \ (\mathbf{const-name} \ \langle \textit{dist} \rangle, \ \textit{NONE}) \rangle$ 
  — To be able to use dist out of the metric-space class.

```

```

class non-decseq-restriction-space =
  uniformity-dist + open-uniformity + restriction- $\sigma$  +
  — We do not assume the restriction sequence to be decseq yet.
  assumes restriction- $\sigma$ -tendsto-zero' :  $\langle \sigma \downarrow \ \textit{TYPE}('a) \longrightarrow 0 \rangle$ 
    and zero-less-restriction- $\sigma'$  [simp] :  $\langle 0 < \sigma \downarrow \ \textit{TYPE}('a) \ n \rangle$ 

    and dist-restriction-is' :  $\langle \textit{dist } x \ y = (\textit{INF } n \in \{n. \ x \downarrow \ n = \ y \downarrow \ n\}).$ 
   $\sigma \downarrow \ \textit{TYPE}('a) \ n \rangle$ 
begin

```

```

sublocale NonDecseqRestrictionSpace  $\langle (\downarrow) \rangle \ \langle (=) \rangle \ \langle \textit{UNIV} \rangle \ \langle \sigma \downarrow \ \textit{TYPE}('a) \rangle \ \textit{dist}$ 
  by unfold-locales (simp-all add: restriction- $\sigma$ -tendsto-zero' dist-restriction-is')

```

end

setup $\langle \text{Sign.add-const-constraint } (\mathbf{const-name} \ \langle \text{dist} \rangle, \text{SOME } \mathbf{typ} \ \langle 'a \rangle$
 $:: \text{metric-space} \Rightarrow 'a \Rightarrow \text{real} \rangle \rangle$
— Only allow *dist* in class *metric-space* (back to normal).

We hide duplicated facts $\sigma_{\downarrow} \text{TYPE}('a) \longrightarrow 0$

$0 < \sigma_{\downarrow} \text{TYPE}('a) \ ?n$

$\text{dist} \ ?x \ ?y = \text{Inf } (\text{restriction-}\sigma\text{-related-set } \ ?x \ ?y).$

hide-fact *restriction- σ -tendsto-zero'* *zero-less-restriction- σ'* *dist-restriction-is'*

context *non-decseq-restriction-space* **begin**

subclass *ultrametric-space*

proof *unfold-locales*

show $\langle \text{dist } x \ y = 0 \iff x = y \rangle$ **for** $x \ y$
by (*simp add: restriction-dist-eq-0-iff-related*)

have *dist-commute* : $\langle \text{dist } x \ y = \text{dist } y \ x \rangle$ **for** $x \ y$
by (*simp add: dist-restriction-is*) *metis*

show $\langle \text{dist } x \ y \leq \max (\text{dist } x \ z) (\text{dist } y \ z) \rangle$ **for** $x \ y \ z$

proof —

consider $\langle x \neq y \rangle$ **and** $\langle y \neq z \rangle$ **and** $\langle x \neq z \rangle$ | $\langle x = y \vee y = z \vee x = z \rangle$ **by** *blast*

thus $\langle \text{dist } x \ y \leq \max (\text{dist } x \ z) (\text{dist } y \ z) \rangle$

proof *cases*

assume $\langle x \neq y \rangle$ **and** $\langle y \neq z \rangle$ **and** $\langle x \neq z \rangle$

from *this(1)[THEN not-related-imp-dist-restriction-le-some-restriction- σ]*

this(2, 3)[THEN not-related-imp-dist-restriction-is-some-restriction- σ]

obtain $l \ m \ n$

where $*$: $\langle \text{dist } x \ y \leq \sigma_{\downarrow} \text{TYPE}('a) \ l \rangle$ $\langle x \downarrow \text{Suc } l \neq y \downarrow \text{Suc } l \rangle$

and $**$: $\langle \text{dist } y \ z = \sigma_{\downarrow} \text{TYPE}('a) \ m \rangle$ $\langle \forall k \leq m. y \downarrow k = z \downarrow k \rangle$

and $***$: $\langle \text{dist } x \ z = \sigma_{\downarrow} \text{TYPE}('a) \ n \rangle$ $\langle \forall k \leq n. x \downarrow k = z \downarrow k \rangle$

by *blast*

have $\langle \min n \ m \leq l \rangle$

proof (*rule ccontr*)

assume $\langle \neg \min n \ m \leq l \rangle$

hence $\langle \text{Suc } l \leq n \rangle$ **and** $\langle \text{Suc } l \leq m \rangle$ **by** *simp-all*

with $** (2)$ $*** (2)$ *Suc-le-lessD*

have $\langle x \downarrow \text{Suc } l = z \downarrow \text{Suc } l \rangle$ $\langle y \downarrow \text{Suc } l = z \downarrow \text{Suc } l \rangle$ **by** *simp-all*

hence $\langle x \downarrow \text{Suc } l = y \downarrow \text{Suc } l \rangle$ **by** *simp*

with $\langle x \downarrow \text{Suc } l \neq y \downarrow \text{Suc } l \rangle$ **show** *False* ..

qed

have $\langle \text{dist } x \ y \leq \sigma_{\downarrow} \text{TYPE}('a) (\min n \ m) \rangle$

by (*rule restriction-space-Inf-properties(1)*)

```

      (simp add: **(2) ***(2))
    also have ⟨... ≤ max (dist x z) (dist y z)⟩
      unfolding **(1) ***(1) by linarith
    finally show ⟨dist x y ≤ max (dist x z) (dist y z)⟩ .
  next
    show ⟨x = y ∨ y = z ∨ x = z ⟹ dist x y ≤ max (dist x z) (dist
y z)⟩
      by (elim disjE, simp-all add: dist-commute)
        (metis not-related-imp-dist-restriction-is-some-restriction-σ
restriction-dist-eq-0-iff-related zero-le-restriction-σ dual-order.refl)
    qed
  qed
qed

end

context non-decseq-restriction-space begin

lemma restriction-tendsto-self: ⟨(λn. x ↓ n) ⟶ x⟩
proof -
  have ⟨(λn. dist (x ↓ n) x) ⟶ 0⟩
  proof (rule real-tendsto-sandwich[OF - ])
    show ⟨∀F n in sequentially. 0 ≤ dist (x ↓ n) x⟩ by simp
  next
    show ⟨∀F n in sequentially. dist (x ↓ n) x ≤ σ↓ TYPE('a) n⟩
    by (auto intro: eventually-sequentiallyI[OF restriction-space-Inf-properties(1)])
  next
    show ⟨(λn. 0) ⟶ 0⟩ by simp
  next
    show ⟨σ↓ TYPE('a) ⟶ 0⟩ by (simp add: restriction-σ-tendsto-zero)
  qed
  thus ⟨(λn. x ↓ n) ⟶ x⟩
  by (subst tendsto-iff[of ⟨(λn. x ↓ n)⟩], subst eventually-sequentially)
    (simp add: LIMSEQ-iff)
qed

end

class decseq-restriction-space = non-decseq-restriction-space +
  assumes decseq-restriction-σ' : ⟨decseq (σ↓ TYPE('a))⟩
begin

sublocale DecseqRestrictionSpace ⟨(↓)⟩ ⟨(=)⟩ ⟨UNIV :: 'a set⟩
  ⟨restriction-σ TYPE('a)⟩ dist

```

by *unfold-locales (simp add: decseq-restriction-σ')*

— Removing $x \in M$ and $y \in M$.

lemmas *dist-restriction-is-bis-simplified* = *dist-restriction-is-bis[simplified]*
and *not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq-simplified*
 =
not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq[simplified]

end

We hide duplicated fact *antimono-on UNIV (σ_↓ TYPE(?'a))*.

hide-fact *decseq-restriction-σ'*

class *strict-decseq-restriction-space* = *non-decseq-restriction-space* +
assumes *strict-decseq-restriction-σ* : $\langle \text{strict-decseq } (\sigma_{\downarrow} \text{ TYPE}'a) \rangle$
begin

subclass *decseq-restriction-space*
 by *unfold-locales*
 (fact *strict-decseq-imp-decseq[OF strict-decseq-restriction-σ]*)

end

Generic Properties

lemma (**in** *metric-space*) *dist-sequences-tendsto-zero-imp-tendsto-iff* :
 $\langle (\lambda n. \text{dist } (\sigma \ n) (\psi \ n)) \longrightarrow 0 \implies \sigma \longrightarrow \Sigma \longleftrightarrow \psi \longrightarrow \Sigma \rangle$

proof (*rule iffI*)

show $\langle \psi \longrightarrow \Sigma \rangle$ **if** $\langle \sigma \longrightarrow \Sigma \rangle$ **and** $\langle (\lambda n. \text{dist } (\sigma \ n) (\psi \ n)) \longrightarrow 0 \rangle$ **for** $\sigma \ \psi$

proof –

from *that* **have** $*$: $\langle (\lambda n. \text{dist } (\sigma \ n) (\psi \ n) + \text{dist } (\sigma \ n) \Sigma) \longrightarrow 0 \rangle$

unfolding *LIMSEQ-dist-iff* **using** *tendsto-add-zero* **by** *blast*

have $\langle (\lambda n. \text{dist } (\psi \ n) \Sigma) \longrightarrow 0 \rangle$

by (*rule real-tendsto-sandwich*)

[*of* $\langle \lambda n. 0 \rangle$ $\langle \lambda n. \text{dist } (\psi \ n) \Sigma \rangle$ - $\langle \lambda n. \text{dist } (\sigma \ n) (\psi \ n) + \text{dist } (\sigma \ n) \Sigma \rangle$]

(*simp-all add: * dist-triangle3*)

with *LIMSEQ-dist-iff* **show** $\langle \psi \longrightarrow \Sigma \rangle$ **by** *blast*

qed

thus $\langle (\lambda n. \text{dist } (\sigma \ n) (\psi \ n)) \longrightarrow 0 \implies \psi \longrightarrow \Sigma \implies \sigma \longrightarrow \Sigma \rangle$

by (*simp add: dist-commute*)

qed

lemma (in *non-decseq-restriction-space*) *restricted-sequence-tendsto-iff*
:
$$\langle (\lambda n. \sigma \ n \downarrow n) \longrightarrow \Sigma \longleftrightarrow \sigma \longrightarrow \Sigma \rangle$$
proof –
$$\text{have } \langle (\lambda n. \text{dist } (\sigma \ n \downarrow n) (\sigma \ n)) \longrightarrow 0 \rangle$$
proof (*unfold metric-space-class.LIMSEQ-def, intro allI impI*)
$$\text{fix } \varepsilon :: \text{real assume } \langle 0 < \varepsilon \rangle$$
from *restriction- σ -tendsto-zero*[*unfolded metric-space-class.LIMSEQ-def, rule-format, OF $\langle 0 < \varepsilon \rangle$*]
$$\text{obtain } no \text{ where } \langle \forall n \geq no. \sigma_{\downarrow} \text{ TYPE}('a) \ n < \varepsilon \rangle$$
by (*auto simp add: dist-real-def*)
$$\text{have } \langle no \leq n \implies \text{dist } (\text{dist } (\sigma \ n \downarrow n) (\sigma \ n)) \ 0 < \varepsilon \rangle \text{ for } n$$
by (*simp, rule le-less-trans*[*OF restriction-space-Inf-properties(1)*][*of $\langle \sigma_{\downarrow} \text{ TYPE}('a) \ n \rangle$*]])
$$\text{(simp, simp add: } \langle \forall n \geq no. \sigma_{\downarrow} \text{ TYPE}('a) \ n < \varepsilon \rangle \text{)}$$
thus $\langle \exists no. \forall n \geq no. \text{dist } (\text{dist } (\sigma \ n \downarrow n) (\sigma \ n)) \ 0 < \varepsilon \rangle$ **by blast**
qed
thus $\langle (\lambda n. \sigma \ n \downarrow n) \longrightarrow \Sigma \longleftrightarrow \sigma \longrightarrow \Sigma \rangle$
by (*simp add: dist-sequences-tendsto-zero-imp-tendsto-iff*)
qed

lemma (in *non-decseq-restriction-space*) *Cauchy-restriction-chain* :
 $\langle \text{Cauchy } \sigma \rangle$ **if** $\langle \text{chain}_{\downarrow} \sigma \rangle$
proof (*rule metric-CauchyI*)
$$\text{fix } \varepsilon :: \text{real assume } \langle 0 < \varepsilon \rangle$$
from *LIMSEQ-D*[*OF restriction- σ -tendsto-zero $\langle 0 < \varepsilon \rangle$, simplified*]
$$\text{obtain } M \text{ where } \langle \sigma_{\downarrow} \text{ TYPE}('a) \ M < \varepsilon \rangle \text{ by blast}$$
moreover $\text{have } \langle M \leq m \implies M \leq n \implies \text{dist } (\sigma \ m) (\sigma \ n) \leq \sigma_{\downarrow} \text{ TYPE}('a) \ M \rangle$ **for** $m \ n$
by (*rule restriction-space-Inf-properties(1), simp add: image-iff*)
$$\text{(metis restriction-chain-def-ter } \langle \text{chain}_{\downarrow} \sigma \rangle \text{)}$$
ultimately $\text{have } \langle \forall m \geq M. \forall n \geq M. \text{dist } (\sigma \ m) (\sigma \ n) < \varepsilon \rangle$
by (*meson dual-order.strict-trans2*)
thus $\langle \exists M. \forall m \geq M. \forall n \geq M. \text{dist } (\sigma \ m) (\sigma \ n) < \varepsilon \rangle$..
qed

lemma (in *non-decseq-restriction-space*) *restriction-tendsto-imp-tendsto*
:
$$\langle \sigma \longrightarrow \Sigma \rangle \text{ if } \langle \sigma \dashrightarrow \Sigma \rangle$$
proof (*rule metric-LIMSEQ-I*)
$$\text{fix } \varepsilon :: \text{real assume } \langle 0 < \varepsilon \rangle$$
with *restriction- σ -tendsto-zero*[*THEN LIMSEQ-D*]
$$\text{obtain } n0 \text{ where } \langle \forall k \geq n0. \sigma_{\downarrow} \text{ TYPE}('a) \ k < \varepsilon \rangle \text{ by auto}$$

hence $\langle \sigma \downarrow \text{TYPE}('a) \ n0 < \varepsilon \rangle$ **by** *simp*
from *restriction-tendstoD[OF $\langle \sigma \ -\downarrow \rightarrow \Sigma \rangle$]*
obtain $n1$ **where** $\langle \forall k \geq n1. \Sigma \downarrow n0 = \sigma \ k \downarrow n0 \rangle ..$
hence $\langle \forall k \geq n1. \text{dist}(\sigma \ k) \Sigma \leq \sigma \downarrow \text{TYPE}('a) \ n0 \rangle$
by (*simp add: restriction-space-Inf-properties(1)*)
with $\langle \sigma \downarrow \text{TYPE}('a) \ n0 < \varepsilon \rangle$
have $\langle \forall k \geq n1. \text{dist}(\sigma \ k) \Sigma < \varepsilon \rangle$ **by** (*meson order-le-less-trans*)
thus $\langle \exists n1. \forall k \geq n1. \text{dist}(\sigma \ k) \Sigma < \varepsilon \rangle ..$
qed

In Decseq Restriction Space

context *decseq-restriction-space* **begin**

lemma *le-dist-to-restriction-eqE* :

obtains k **where** $\langle n \leq k \rangle \langle \bigwedge x \ y :: 'a. \text{dist} \ x \ y \leq \sigma \downarrow \text{TYPE}('a) \ k \implies x \downarrow n = y \downarrow n \rangle$

proof –

have $\langle \exists k \geq n. \forall x \ y :: 'a. \text{dist} \ x \ y \leq \sigma \downarrow \text{TYPE}('a) \ k \longrightarrow x \downarrow n = y \downarrow n \rangle$

proof (*rule ccontr*)

assume $\langle \neg (\exists k \geq n. \forall x \ y :: 'a. \text{dist} \ x \ y \leq \sigma \downarrow \text{TYPE}('a) \ k \longrightarrow x \downarrow n = y \downarrow n) \rangle$

hence $\langle \forall k \geq n. \exists x \ y :: 'a. \text{dist} \ x \ y \leq \sigma \downarrow \text{TYPE}('a) \ k \wedge x \downarrow n \neq y \downarrow n \rangle$ **by** *simp*

then obtain $X \ Y :: \langle \text{nat} \Rightarrow 'a \rangle$

where $*$: $\langle \forall k \geq n. \text{dist} \ (X \ k) \ (Y \ k) \leq \sigma \downarrow \text{TYPE}('a) \ k \rangle$

$\langle \forall k \geq n. X \ k \downarrow n \neq Y \ k \downarrow n \rangle$ **by** *metis*

moreover obtain $n0$ **where** $\langle \forall k \geq n0. \sigma \downarrow \text{TYPE}('a) \ k < \sigma \downarrow \text{TYPE}('a) \ n \rangle$

by (*metis LIMSEQ-D zero-less-restriction- σ abs-of-nonneg diff-zero restriction- σ -tendsto-zero real-norm-def zero-le-restriction- σ*)

ultimately have $\langle \forall k \geq n+n0. \text{dist} \ (X \ k) \ (Y \ k) < \sigma \downarrow \text{TYPE}('a) \ n \rangle$

by (*metis dual-order.strict-trans2 add-leE*)

moreover from $*$ (2) **have** $\langle \forall k \geq n. \sigma \downarrow \text{TYPE}('a) \ n \leq \text{dist} \ (X \ k) \ (Y \ k) \rangle$

by (*simp add: dist-restriction-is-bis*)

(*metis (full-types) decseq-restriction- σ [THEN decseqD] linorder-linear*

not-eq-imp-dist-restriction-is-restriction- σ -Sup-restriction-eq-simplified(2))

ultimately show *False* **by** (*metis le-add1 linorder-not-le order-refl*)

qed

thus $\langle (\bigwedge k. \llbracket n \leq k; \bigwedge x \ y :: 'a. \text{dist} \ x \ y \leq \sigma \downarrow \text{TYPE}('a) \ k \implies x \downarrow n = y \downarrow n \rrbracket$

$\implies \text{thesis} \rangle \implies \text{thesis}$ **by** *blast*

qed

theorem *tendsto-iff-restriction-tendsto* : $\langle \sigma \longrightarrow \Sigma \longleftrightarrow \sigma \ -\downarrow \rightarrow \Sigma \rangle$

proof (*rule iffI*)

show $\langle \sigma \dashv \rightarrow \Sigma \implies \sigma \longrightarrow \Sigma \rangle$ **by** (*fact restriction-tendsto-imp-tendsto*)
next
show $\langle \sigma \dashv \rightarrow \Sigma \rangle$ **if** $\langle \sigma \longrightarrow \Sigma \rangle$
proof (*rule restriction-tendstoI*)
fix n
obtain $n0$ **where** $\langle \text{dist } x \ y \leq \sigma \downarrow \text{TYPE}('a) \ n0 \implies x \downarrow n = y \downarrow n \rangle$
for $x \ y :: 'a$
by (*metis le-dist-to-restriction-eqE*)
moreover from *metric-LIMSEQ-D*[*OF* $\langle \sigma \longrightarrow \Sigma \rangle$ *zero-less-restriction- σ*]
obtain $n1$ **where** $\langle \forall k \geq n1. \text{dist } (\sigma \ k) \ \Sigma < \sigma \downarrow \text{TYPE}('a) \ n0 \rangle ..$
ultimately have $\langle \forall k \geq n1. \Sigma \downarrow n = \sigma \ k \downarrow n \rangle$ **by** (*metis dual-order.order-iff-strict*)
thus $\langle \exists n1. \forall k \geq n1. \Sigma \downarrow n = \sigma \ k \downarrow n \rangle ..$
qed
qed

corollary *convergent-iff-restriction-convergent* : $\langle \text{convergent } \sigma \longleftrightarrow \text{convergent}_\downarrow \sigma \rangle$
by (*simp add: convergent-def restriction-convergent-def tendsto-iff-restriction-tendsto*)

theorem *complete-iff-restriction-complete* :

$\langle (\forall \sigma. \text{Cauchy } \sigma \longrightarrow \text{convergent } \sigma) \longleftrightarrow (\forall \sigma. \text{chain}_\downarrow \sigma \longrightarrow \text{convergent}_\downarrow \sigma) \rangle$

— The following result shows that we have not lost anything with our definitions of convergence, completeness, etc. specific to restriction spaces.

proof (*intro iffI impI allI*)

fix σ **assume** *hyp* : $\langle \forall \sigma. \text{Cauchy } \sigma \longrightarrow \text{convergent } \sigma \rangle$ **and** $\langle \text{chain}_\downarrow \sigma \rangle$

from *Cauchy-restriction-chain* $\langle \text{chain}_\downarrow \sigma \rangle$ **have** $\langle \text{Cauchy } \sigma \rangle$ **by** *blast*

hence $\langle \text{convergent } \sigma \rangle$ **by** (*simp add: hyp*)

thus $\langle \text{convergent}_\downarrow \sigma \rangle$ **by** (*simp add: convergent-iff-restriction-convergent*)

next

fix σ **assume** *hyp* : $\langle \forall \sigma. \text{chain}_\downarrow \sigma \longrightarrow \text{convergent}_\downarrow \sigma \rangle$ **and** $\langle \text{Cauchy } \sigma \rangle$

from $\langle \text{Cauchy } \sigma \rangle$ **have** $*$: $\langle \forall n. \exists k. \forall l \geq k. \sigma \ l \downarrow n = \sigma \ k \downarrow n \rangle$

by (*metis (mono-tags, opaque-lifting) Cauchy-altdef2 dual-order.order-iff-strict*)

le-dist-to-restriction-eqE zero-less-restriction- σ)

define f **where** $\langle f \equiv \text{rec-nat}$

$(\text{LEAST } k. \forall l \geq k. \sigma \ l \downarrow 0 = \sigma \ k \downarrow 0)$

$(\lambda n \ k. \text{LEAST } l. k < l \wedge (\forall m \geq l. \sigma \ m \downarrow \text{Suc } n = \sigma \ l$

$\downarrow \text{Suc } n)) \rangle$

have *f-Suc-def* : $\langle f (\text{Suc } n) = (\text{LEAST } l. f \ n < l \wedge (\forall m \geq l. \sigma \ m \downarrow \text{Suc } n = \sigma \ l \downarrow \text{Suc } n)) \rangle$

(is $\langle f (\text{Suc } n) = \text{Least } (?f\text{-Suc } n) \rangle$) **for** n **by** (*simp add: f-def*)

from $*$ **have** $**$: $\langle \exists k > f \ n. \forall m \geq k. \sigma \ m \downarrow \text{Suc } n = \sigma \ k \downarrow \text{Suc } n \rangle$ **for** n

by (*metis dual-order.trans lessI linorder-not-le order-le-less*)

```

have ⟨strict-mono f⟩
proof (unfold strict-mono-Suc-iff, rule allI)
  show ⟨f n < f (Suc n)⟩ for n
    by (fact LeastI-ex[of ⟨?f-Suc n⟩, folded f-Suc-def, OF **, THEN
conjunct1])
qed

have ⟨chain↓ ( $\lambda n. (\sigma \circ f) n \downarrow n$ )⟩ — key point
proof (rule restriction-chainI)
  fix n
  have ⟨ $\sigma (f (Suc\ n)) \downarrow n = \sigma (f\ n) \downarrow n$ ⟩
proof (cases n)
  show ⟨n = 0  $\implies \sigma (f (Suc\ n)) \downarrow n = \sigma (f\ n) \downarrow n$ ⟩ by simp
next
  show ⟨n = Suc nat  $\implies \sigma (f (Suc\ n)) \downarrow n = \sigma (f\ n) \downarrow n$ ⟩ for nat
proof (clarify, intro LeastI-ex[of ⟨?f-Suc nat⟩,
  folded f-Suc-def, THEN conjunct2, rule-format, OF **])
  show ⟨n = Suc nat  $\implies f (Suc\ nat) \leq f (Suc (Suc\ nat))$ ⟩
    by (simp add: ⟨strict-mono f⟩ strict-mono-less-eq)
qed
qed
thus ⟨ $(\sigma \circ f) (Suc\ n) \downarrow Suc\ n \downarrow n = (\sigma \circ f) n \downarrow n$ ⟩ by simp
qed
with hyp have ⟨convergent↓ ( $\lambda n. (\sigma \circ f) n \downarrow n$ )⟩ by simp
hence ⟨convergent↓ ( $\lambda n. (\sigma \circ f) n$ )⟩
by (simp add: restriction-convergent-restricted-iff-restriction-convergent)
hence ⟨convergent ( $\lambda n. (\sigma \circ f) n$ )⟩
by (simp add: convergent-iff-restriction-convergent)
with Cauchy-converges-subseq[OF ⟨Cauchy  $\sigma$ ⟩ ⟨strict-mono f⟩]
show ⟨convergent  $\sigma$ ⟩ unfolding convergent-def by blast
qed

end

The following classes will be useful later.

class complete-decseq-restriction-space = decseq-restriction-space +
  assumes restriction-chain-imp-restriction-convergent' : ⟨chain↓  $\sigma$ 
 $\implies$  convergent↓  $\sigma$ ⟩
begin

subclass complete-restriction-space
  by unfold-locales (fact restriction-chain-imp-restriction-convergent')

subclass complete-ultrametric-space
proof (unfold-locales)
from complete-iff-restriction-complete restriction-chain-imp-restriction-convergent
show ⟨Cauchy  $\sigma \implies$  convergent  $\sigma$ ⟩ for  $\sigma$  by blast
qed

```

end

We hide duplicated fact $chain_{\downarrow} \ ?\sigma \implies convergent_{\downarrow} \ ?\sigma$.

hide-fact *restriction-chain-imp-restriction-convergent'*

class *complete-strict-decseq-restriction-space* = *strict-decseq-restriction-space* +

assumes *restriction-chain-imp-restriction-convergent'* : $\langle chain_{\downarrow} \ \sigma \implies convergent_{\downarrow} \ \sigma \rangle$

begin

subclass *complete-decseq-restriction-space*

by (*unfold-locales*) (*fact restriction-chain-imp-restriction-convergent'*)

end

We hide duplicated fact $chain_{\downarrow} \ ?\sigma \implies convergent_{\downarrow} \ ?\sigma$.

hide-fact *restriction-chain-imp-restriction-convergent'*

class *restriction- δ* = *restriction- σ* +

fixes *restriction- δ* :: $\langle 'a \ \textit{itself} \ \Rightarrow \ \textit{real} \rangle \ (\delta_{\downarrow})$

assumes *zero-less-restriction- δ* [*simp*] : $\langle 0 < \delta_{\downarrow} \ \textit{TYPE}('a) \rangle$

and *restriction- δ -less-one* [*simp*] : $\langle \delta_{\downarrow} \ \textit{TYPE}('a) < 1 \rangle$

begin

lemma *zero-le-restriction- δ* [*simp*] : $\langle 0 \leq \delta_{\downarrow} \ \textit{TYPE}('a) \rangle$

and *restriction- δ -le-one* [*simp*] : $\langle \delta_{\downarrow} \ \textit{TYPE}('a) \leq 1 \rangle$

and *zero-le-pow-restriction- δ* [*simp*] : $\langle 0 \leq \delta_{\downarrow} \ \textit{TYPE}('a) \wedge n \rangle$

by (*simp-all add: order-less-imp-le*)

lemma *pow-restriction- δ -le-one* [*simp*] : $\langle \delta_{\downarrow} \ \textit{TYPE}('a) \wedge n \leq 1 \rangle$

by (*simp add: power-le-one*)

lemma *pow-restriction- δ -less-one* [*simp*] : $\langle n \neq 0 \implies \delta_{\downarrow} \ \textit{TYPE}('a) \wedge n < 1 \rangle$

by (*metis restriction- δ -less-one zero-less-restriction- δ not-gr-zero power-0 power-strict-decreasing*)

end

setup $\langle \textit{Sign.add-const-constraint} \ (\mathbf{const-name} \ \langle \textit{dist} \rangle, \ \textit{NONE}) \rangle$

 — To be able to use *dist* out of the *metric-space* class.

class *at-least-geometric-restriction-space* =

uniformity-dist + *open-uniformity* + *restriction- δ* +

```

assumes zero-less-restriction- $\sigma'$  :  $\langle 0 < \sigma_{\downarrow} \text{TYPE}('a) \ n \rangle$ 
and restriction- $\sigma$ -le :
 $\langle \sigma_{\downarrow} \text{TYPE}('a) \ (\text{Suc } n) \leq \delta_{\downarrow} \text{TYPE}('a) * \sigma_{\downarrow} \text{TYPE}('a) \ n \rangle$ 
and dist-restriction-is' :
 $\langle \text{dist } x \ y = (\text{INF } n \in \{n. \ x \downarrow n = y \downarrow n\}. \ \sigma_{\downarrow} \text{TYPE}('a) \ n) \rangle$ 

setup  $\langle \text{Sign.add-const-constraint } (\text{const-name } \langle \text{dist} \rangle, \text{ SOME } \text{typ } \langle 'a \rangle$ 
 $:: \text{metric-space} \Rightarrow 'a \Rightarrow \text{real} \rangle \rangle$ 
— Only allow dist in class metric-space (back to normal).

context at-least-geometric-restriction-space begin

lemma restriction- $\sigma$ -le-restriction- $\sigma$ -times-pow-restriction- $\delta$  :
 $\langle \sigma_{\downarrow} \text{TYPE}('a) \ (n + k) \leq \sigma_{\downarrow} \text{TYPE}('a) \ n * \delta_{\downarrow} \text{TYPE}('a) \ ^k \rangle$ 
by (induct k, simp-all)
(metis dual-order.trans restriction- $\sigma$ -le zero-less-restriction- $\delta$ 
mult.left-commute mult.le-cancel-left-pos)

lemma restriction- $\sigma$ -le-pow-restriction- $\delta$  :
 $\langle \sigma_{\downarrow} \text{TYPE}('a) \ n \leq \sigma_{\downarrow} \text{TYPE}('a) \ 0 * \delta_{\downarrow} \text{TYPE}('a) \ ^n \rangle$ 
by (metis add-0 restriction- $\sigma$ -le-restriction- $\sigma$ -times-pow-restriction- $\delta$ )

subclass strict-decseq-restriction-space
proof unfold-locales
have  $\langle \forall_F \ n \text{ in sequentially. } 0 \leq \sigma_{\downarrow} \text{TYPE}('a) \ n \rangle$ 
by (simp add: zero-less-restriction- $\sigma'$  order-less-imp-le)
moreover have  $\langle \forall_F \ n \text{ in sequentially. } \sigma_{\downarrow} \text{TYPE}('a) \ n$ 
 $\leq \sigma_{\downarrow} \text{TYPE}('a) \ 0 * \delta_{\downarrow} \text{TYPE}('a) \ ^n \rangle$ 
by (simp add: restriction- $\sigma$ -le-pow-restriction- $\delta$ )
moreover have  $\langle (\lambda n. \ 0) \longrightarrow 0 \rangle$  by simp
moreover have  $\langle (\lambda n. \ \sigma_{\downarrow} \text{TYPE}('a) \ 0 * \delta_{\downarrow} \text{TYPE}('a) \ ^n) \longrightarrow$ 
 $0 \rangle$ 
by (simp add: LIMSEQ-abs-realpow-zero2 abs-of-pos tendsto-mult-right-zero)
ultimately show  $\langle \sigma_{\downarrow} \text{TYPE}('a) \longrightarrow 0 \rangle$  by (fact real-tendsto-sandwich)
next
show  $\langle 0 < \sigma_{\downarrow} \text{TYPE}('a) \ n \rangle$  for n
by (fact zero-less-restriction- $\sigma'$ )
next
show  $\langle \text{dist } x \ y = \text{Inf } (\sigma_{\downarrow} \text{TYPE}('a) \ \{n. \ x \downarrow n = y \downarrow n\}) \rangle$  for x y
by (fact dist-restriction-is')
next
show  $\langle \text{strict-decseq } (\sigma_{\downarrow} \text{TYPE}('a)) \rangle$ 
by (rule strict-decseqI, rule le-less-trans[OF restriction- $\sigma$ -le])
(simp add: zero-less-restriction- $\sigma'$ )
qed

```

lemma $\langle 0 < \delta_{\downarrow} \text{TYPE}('a) \wedge n \rangle$ **by** *simp*

end

We hide duplicated facts $0 < \sigma_{\downarrow} \text{TYPE}('a) \ ?n$
 $\text{dist } ?x \ ?y = \text{Inf } (\text{restriction-}\sigma\text{-related-set } ?x \ ?y)$.

hide-fact *zero-less-restriction- σ' dist-restriction-is'*

class *complete-at-least-geometric-restriction-space* = *at-least-geometric-restriction-space*
 +

assumes *restriction-chain-imp-restriction-convergent'* : $\langle \text{chain}_{\downarrow} \sigma$
 $\implies \text{convergent}_{\downarrow} \sigma \rangle$

begin

subclass *complete-strict-decseq-restriction-space*

by *unfold-locales* (*fact restriction-chain-imp-restriction-convergent'*)

end

We hide duplicated fact $\text{chain}_{\downarrow} \ ?\sigma \implies \text{convergent}_{\downarrow} \ ?\sigma$.

hide-fact *restriction-chain-imp-restriction-convergent'*

setup $\langle \text{Sign.add-const-constraint } (\mathbf{const-name} \ \langle \text{dist} \rangle, \text{NONE}) \rangle$

— To be able to use *dist* out of the *metric-space* class.

class *geometric-restriction-space* = *uniformity-dist* + *open-uniformity*
 + *restriction- δ* +

assumes *restriction- σ -is* : $\langle \sigma_{\downarrow} \text{TYPE}('a) \ n = \delta_{\downarrow} \text{TYPE}('a) \wedge n \rangle$

and *dist-restriction-is'* : $\langle \text{dist } x \ y = (\text{INF } n \in \{n. x \downarrow n = y \downarrow n\}).$
 $\sigma_{\downarrow} \text{TYPE}('a) \ n \rangle$

begin

This is what “restriction space” usually mean in the literature.
 The previous classes are generalizations of this concept (even this
 one is a generalization, since we usually have $\delta_{\downarrow} \text{TYPE}('a) = 1$
 / 2).

subclass *at-least-geometric-restriction-space*

proof *unfold-locales*

show $\langle 0 < \sigma_{\downarrow} \text{TYPE}('a) \ n \rangle$ **for** n **by** (*simp add: restriction- σ -is*)

next

```

show  $\langle \sigma_{\downarrow} \text{TYPE}('a) (\text{Suc } n) \leq \delta_{\downarrow} \text{TYPE}('a) * \sigma_{\downarrow} \text{TYPE}('a) n \rangle$  for
 $n$ 
  by (simp add: restriction- $\sigma$ -is)
next
show  $\langle \text{dist } x y = \text{Inf } (\sigma_{\downarrow} \text{TYPE}('a) \{n. x \downarrow n = y \downarrow n\}) \rangle$  for  $x y$ 
  by (fact dist-restriction-is')
qed

```

```

lemma  $\langle 0 < \delta_{\downarrow} \text{TYPE}('a) \wedge n \rangle$  by simp

```

```

end

```

```

setup  $\langle \text{Sign.add-const-constraint } (\text{const-name } \langle \text{dist} \rangle, \text{SOME } \text{typ } \langle 'a$ 
 $:: \text{metric-space} \Rightarrow 'a \Rightarrow \text{real} \rangle) \rangle$ 
  — Only allow dist in class metric-space (back to normal).

```

We hide duplicated fact $\text{dist } ?x ?y = \text{Inf } (\text{restriction-}\sigma\text{-related-set } ?x ?y)$.

```

hide-fact dist-restriction-is'

```

```

class complete-geometric-restriction-space = geometric-restriction-space
+
  assumes restriction-chain-imp-restriction-convergent' :  $\langle \text{chain}_{\downarrow} \sigma$ 
 $\Longrightarrow \text{convergent}_{\downarrow} \sigma \rangle$ 
begin

```

```

subclass complete-at-least-geometric-restriction-space
  by (unfold-locales) (fact restriction-chain-imp-restriction-convergent')

```

```

end

```

We hide duplicated fact $\text{chain}_{\downarrow} ?\sigma \Longrightarrow \text{convergent}_{\downarrow} ?\sigma$.

```

hide-fact restriction-chain-imp-restriction-convergent'

```

```

theorem geometric-restriction-space-completeI :  $\langle \text{convergent } \sigma \rangle$ 
if  $\langle \bigwedge \sigma :: \text{nat} \Rightarrow 'a. \text{restriction-chain } \sigma \Longrightarrow \exists \Sigma. \forall n. \Sigma \downarrow n = \sigma n \rangle$ 
and  $\langle \text{Cauchy } \sigma \rangle$  for  $\sigma :: \langle \text{nat} \Rightarrow 'a :: \text{geometric-restriction-space} \rangle$ 
by (metis complete-iff-restriction-complete convergent-def
convergent-iff-restriction-convergent ext restriction-tendsto-self
that)

```

— Because *cball* is not defined in *metric-space*.

```

lemma (in non-decseq-restriction-space) restriction-cball-subset-cball
:

```

$\langle \sigma_{\downarrow} \text{TYPE}('a) \ n \leq r \implies \mathcal{B}_{\downarrow}(\Sigma, n) \subseteq \{x. \text{dist } \Sigma \ x \leq r\} \rangle$
by (*simp add: subset-iff restriction-cball-mem-iff*)
(simp add: dual-order.trans restriction-space-Inf-properties(1))

corollary *restriction-cball-subset-cball-bis* :
 $\langle \sigma_{\downarrow} \text{TYPE}('a) \ n \leq r \implies \mathcal{B}_{\downarrow}(\Sigma, n) \subseteq \text{cball } \Sigma \ r \rangle$
for $\Sigma :: \langle 'a :: \text{non-decseq-restriction-space} \rangle$
unfolding *cball-def* **by** (*fact restriction-cball-subset-cball*)

lemma (*in non-decseq-restriction-space*) *restriction-ball-subset-ball* :
 $\langle \sigma_{\downarrow} \text{TYPE}('a) \ n < r \implies \text{restriction-ball } \Sigma \ n \subseteq \{x. \text{dist } \Sigma \ x < r\} \rangle$
by (*simp add: subset-iff restriction-cball-mem-iff*)
(metis (mono-tags, lifting) image-eqI mem-Collect-eq order-le-less-trans restriction-related-pred restriction-space-Inf-properties(1))

corollary *restriction-ball-subset-ball-bis* :
 $\langle \sigma_{\downarrow} \text{TYPE}('a) \ n < r \implies \text{restriction-ball } \Sigma \ n \subseteq \text{ball } \Sigma \ r \rangle$
for $\Sigma :: \langle 'a :: \text{non-decseq-restriction-space} \rangle$
unfolding *ball-def* **by** (*fact restriction-ball-subset-ball*)

lemma (*in strict-decseq-restriction-space*)
restriction-cball-is-cball : $\langle \mathcal{B}_{\downarrow}(\Sigma, n) = \{x. \text{dist } \Sigma \ x \leq \sigma_{\downarrow} \text{TYPE}('a) \ n\} \rangle$
proof (*intro subset-antisym subsetI*)
from *restriction-cball-subset-cball*
show $\langle x \in \mathcal{B}_{\downarrow}(\Sigma, n) \implies x \in \{x. \text{dist } \Sigma \ x \leq \sigma_{\downarrow} \text{TYPE}('a) \ n\} \rangle$ **for**
 x **by** *blast*
next
show $\langle x \in \mathcal{B}_{\downarrow}(\Sigma, n) \rangle$ **if** $\langle x \in \{x. \text{dist } \Sigma \ x \leq \sigma_{\downarrow} \text{TYPE}('a) \ n\} \rangle$ **for** x
proof (*cases* $\langle \Sigma = x \rangle$)
show $\langle \Sigma = x \implies x \in \mathcal{B}_{\downarrow}(\Sigma, n) \rangle$ **by** *simp*
next
assume $\langle \Sigma \neq x \rangle$
with *not-related-imp-dist-restriction-is-some-restriction- σ* **obtain**
 m
where $\langle \text{dist } \Sigma \ x = \sigma_{\downarrow} \text{TYPE}('a) \ m \rangle \langle \forall k \leq m. \Sigma \downarrow k = x \downarrow k \rangle$ **by**
blast
from $\langle x \in \{x. \text{dist } \Sigma \ x \leq \sigma_{\downarrow} \text{TYPE}('a) \ n\} \rangle$
have $\langle \text{dist } \Sigma \ x \leq \sigma_{\downarrow} \text{TYPE}('a) \ n \rangle$ **by** *simp*
with $\langle \text{dist } \Sigma \ x = \sigma_{\downarrow} \text{TYPE}('a) \ m \rangle$ **have** $\langle n \leq m \rangle$
by (*metis linorder-not-less strict-decseq-restriction- σ strict-decseq-def-bis*)
with $\langle \forall k \leq m. \Sigma \downarrow k = x \downarrow k \rangle$ **have** $\langle \Sigma \downarrow n = x \downarrow n \rangle$ **by** *simp*
thus $\langle x \in \mathcal{B}_{\downarrow}(\Sigma, n) \rangle$ **by** (*simp add: restriction-cball-mem-iff*)
qed
qed

```

lemma restriction-cball-is-cball-bis :  $\langle \mathcal{B}_\downarrow(\Sigma, n) = \text{cball } \Sigma (\sigma_\downarrow \text{TYPE}('a) n) \rangle$ 
for  $\Sigma :: \langle 'a :: \text{strict-decseq-restriction-space} \rangle$ 
by (simp add: cball-def restriction-cball-is-cball)

lemma (in strict-decseq-restriction-space)
  restriction-ball-is-ball :  $\langle \text{restriction-ball } \Sigma n = \{x. \text{dist } \Sigma x < \sigma_\downarrow \text{TYPE}('a) n\} \rangle$ 
proof (intro subset-antisym subsetI)
  show  $\langle x \in \text{restriction-ball } \Sigma n \implies x \in \{x. \text{dist } \Sigma x < \sigma_\downarrow \text{TYPE}('a) n\} \rangle$ 
  for  $x$ 
    by (simp add: subset-iff)
    (metis lessI local.restriction-cball-is-cball strict-decseq-restriction- $\sigma$  mem-Collect-eq order-le-less-trans strict-decseq-def-bis)
next
  show  $\langle x \in \text{restriction-ball } \Sigma n \rangle$  if  $\langle x \in \{x. \text{dist } \Sigma x < \sigma_\downarrow \text{TYPE}('a) n\} \rangle$ 
  for  $x$ 
    proof (cases  $\langle \Sigma = x \rangle$ )
      show  $\langle \Sigma = x \implies x \in \text{restriction-ball } \Sigma n \rangle$  by simp
    next
      assume  $\langle \Sigma \neq x \rangle$ 
      with not-related-imp-dist-restriction-is-some-restriction- $\sigma$  obtain
         $m$ 
        where  $\langle \text{dist } \Sigma x = \sigma_\downarrow \text{TYPE}('a) m \rangle \langle \forall k \leq m. \Sigma \downarrow k = x \downarrow k \rangle$  by
          blast
        from  $\langle x \in \{x. \text{dist } \Sigma x < \sigma_\downarrow \text{TYPE}('a) n\} \rangle$ 
        have  $\langle \text{dist } \Sigma x < \sigma_\downarrow \text{TYPE}('a) n \rangle$  by simp
        with  $\langle \text{dist } \Sigma x = \sigma_\downarrow \text{TYPE}('a) m \rangle$  have  $\langle n < m \rangle$ 
        using strict-decseq-restriction- $\sigma$  strict-decseq-def-bis by auto
        with  $\langle \forall k \leq m. \Sigma \downarrow k = x \downarrow k \rangle$  have  $\langle \Sigma \downarrow \text{Suc } n = x \downarrow \text{Suc } n \rangle$  by
          simp
        thus  $\langle x \in \text{restriction-ball } \Sigma n \rangle$  by (simp add: restriction-cball-mem-iff)
      qed
    qed

```

```

lemma restriction-ball-is-ball-bis :  $\langle \text{restriction-ball } \Sigma n = \text{ball } \Sigma (\sigma_\downarrow \text{TYPE}('a) n) \rangle$ 
for  $\Sigma :: \langle 'a :: \text{strict-decseq-restriction-space} \rangle$ 
by (simp add: ball-def restriction-ball-is-ball)

```

```

lemma isCont-iff-restriction-cont-at :  $\langle \text{isCont } f \Sigma \iff \text{restriction-cont-at } f \Sigma \rangle$ 
for  $f :: \langle 'a :: \text{decseq-restriction-space} \implies 'b :: \text{decseq-restriction-space} \rangle$ 
by (unfold restriction-cont-at-def continuous-at-sequentially-comp-def,

```

fold tendsto-iff-restriction-tendsto) simp

lemma (in *strict-decseq-restriction-space*)
open-iff-restriction-open : $\langle \text{open } U \longleftrightarrow \text{open}_\downarrow U \rangle$
proof (*unfold open-dist restriction-open-iff-restriction-cball-characterization,*
intro iffI ballI)
fix Σ **assume** $\langle \forall x \in U. \exists e > 0. \forall y. \text{dist } y \ x < e \longrightarrow y \in U \rangle$ **and** $\langle \Sigma$
 $\in U \rangle$
then obtain e **where** $\langle 0 < e \rangle \langle \forall y. \text{dist } y \ \Sigma < e \longrightarrow y \in U \rangle$ **by**
blast
from $\langle 0 < e \rangle$ **obtain** n **where** $\langle \sigma_\downarrow \text{TYPE}('a) \ n < e \rangle$
by (*metis eventually-at-top-linorder le-refl restriction- σ -tendsto-zero*
order-tendstoD(2))
with $\langle \forall y. \text{dist } y \ \Sigma < e \longrightarrow y \in U \rangle$ *dist-commute restriction-cball-is-cball*
have $\langle \mathcal{B}_\downarrow(\Sigma, n) \subseteq U \rangle$ **by** *auto*
thus $\langle \exists n. \mathcal{B}_\downarrow(\Sigma, n) \subseteq U \rangle$..
next
fix x **assume** $\langle \forall \Sigma \in U. \exists n. \mathcal{B}_\downarrow(\Sigma, n) \subseteq U \rangle \langle x \in U \rangle$
then obtain n **where** $\langle \mathcal{B}_\downarrow(x, n) \subseteq U \rangle$ **by** *blast*
hence $\langle \forall y. \text{dist } y \ x < \sigma_\downarrow \text{TYPE}('a) \ n \longrightarrow y \in U \rangle$
by (*simp add: dist-commute restriction-cball-is-cball subset-iff*)
thus $\langle \exists e > 0. \forall y. \text{dist } y \ x < e \longrightarrow y \in U \rangle$
using *zero-less-restriction- σ'* **by** *blast*
qed

lemma (in *strict-decseq-restriction-space*)
closed-iff-restriction-closed : $\langle \text{closed } U \longleftrightarrow \text{closed}_\downarrow U \rangle$
by (*simp add: closed-open open-iff-restriction-open restriction-open-Compl-iff*)

lemma *continuous-on-iff-restriction-cont-on* :
 $\langle \text{open } U \implies \text{continuous-on } U \ f \longleftrightarrow \text{restriction-cont-on } f \ U \rangle$
for $f :: 'a :: \text{decseq-restriction-space} \Rightarrow 'b :: \text{decseq-restriction-space}$
by (*simp add: restriction-cont-on-def continuous-on-eq-continuous-at*
flip: isCont-iff-restriction-cont-at)

3.2 Equivalence between Lipschitz Map and Restriction shift Map

For a function $f :: 'a \Rightarrow 'b$, it is equivalent to have *restriction-shift-on* $f \ k \ A$ and *lipschitz-with-on* $f \ (\delta_\downarrow \text{TYPE}('b))^k \ A$ when $'a$ is of sort *geometric-restriction-space* and $\sigma_\downarrow \text{TYPE}('b) = \sigma_\downarrow \text{TYPE}('a)$ ($'b$ is then necessarily also of sort *geometric-restriction-space*).

Weaker versions of this result can be established with weaker

assumptions on the sort, this is what we do below.

lemma *restriction-shift-nonneg-on-imp-lipschitz-with-on* :
 ‹lipschitz-with-on f (restriction- δ TYPE('b) \wedge k) A› **if** ‹restriction-shift-on f (int k) A›
and le-restriction- σ : ‹ $\bigwedge n$. restriction- σ TYPE('b) n \leq restriction- σ TYPE('a) n›
for f :: ‹'a :: decseq-restriction-space \Rightarrow 'b :: at-least-geometric-restriction-space›
proof (rule lipschitz-with-onI)
show ‹0 \leq restriction- δ TYPE('b) \wedge k› **by** simp
next
fix x y **assume** ‹x \in A› ‹y \in A› ‹x \neq y› ‹f x \neq f y›
from ‹restriction-shift-on f k A› [THEN restriction-shift-onD, OF ‹x \in A› ‹y \in A›]
have ‹i \in restriction-related-set x y \implies i + k \in restriction-related-set (f x) (f y)› **for** i
by (simp add: nat-int-add)
hence ‹Sup (restriction-related-set x y) + k \in restriction-related-set (f x) (f y)›
using ‹x \neq y› Sup-in-restriction-related-set **by** blast
hence ‹Sup (restriction-related-set x y) + k \leq Sup (restriction-related-set (f x) (f y))›
by (simp add: ‹f x \neq f y› bdd-above-restriction-related-set-iff cSup-upper)
moreover have ‹dist (f x) (f y) = restriction- σ TYPE('b) (Sup (restriction-related-set (f x) (f y)))›
by (simp add: ‹f x \neq f y› dist-restriction-is-bis-simplified)
ultimately have ‹dist (f x) (f y) \leq restriction- σ TYPE('b) (Sup (restriction-related-set x y) + k)›
by (simp add: decseqD decseq-restriction-space-class.decseq-restriction- σ)
hence ‹dist (f x) (f y) \leq restriction- σ TYPE('b) (Sup (restriction-related-set x y) * restriction- δ TYPE('b) \wedge k)›
by (meson order.trans restriction- σ -le-restriction- σ -times-pow-restriction- δ)
also have ‹... \leq restriction- σ TYPE('a) (Sup (restriction-related-set x y) * restriction- δ TYPE('b) \wedge k)›
by (simp add: le-restriction- σ)
finally show ‹dist (f x) (f y) \leq restriction- δ TYPE('b) \wedge k * dist x y›
by (simp add: dist-restriction-is-bis[of x y] ‹x \neq y› mult.commute)
qed

corollary *restriction-shift-nonneg-imp-lipschitz-with* :
 ‹[[restriction-shift f (int k); $\bigwedge n$. restriction- σ TYPE('b) n \leq restriction- σ TYPE('a) n]]
 \implies lipschitz-with f (restriction- δ TYPE('b) \wedge k)›
for f :: ‹'a :: decseq-restriction-space \Rightarrow 'b :: at-least-geometric-restriction-space›
using restriction-shift-def restriction-shift-nonneg-on-imp-lipschitz-with-on
by blast

lemma *lipschitz-with-on-imp-restriction-shift-neg-on* :

$\langle \text{restriction-shift-on } f \text{ } (- \text{ int } k) \text{ } A \rangle$ **if** $\langle \text{lipschitz-with-on } f \text{ } (\text{restriction-}\delta \text{ } \text{TYPE}'(b) \text{ } \text{powi} - \text{ int } k) \text{ } A \rangle$

and $\langle \text{le-restriction-}\sigma : \langle \bigwedge n. \text{restriction-}\sigma \text{ } \text{TYPE}'(a) \text{ } n \leq \text{restriction-}\sigma \text{ } \text{TYPE}'(b) \text{ } n \rangle$

for $f :: \langle 'a :: \text{decseq-restriction-space} \Rightarrow 'b :: \text{at-least-geometric-restriction-space} \rangle$

proof (*rule restriction-shift-onI, goal-cases*)

fix $x \ y \ n$ **assume** $\langle x \in A \rangle \langle y \in A \rangle \langle f \ x \neq f \ y \rangle \langle x \downarrow n = y \downarrow n \rangle$

from $\langle f \ x \neq f \ y \rangle$ **have** $\langle x \neq y \rangle$ **by** *blast*

from $\langle \text{lipschitz-with-on } f \text{ } (\text{restriction-}\delta \text{ } \text{TYPE}'(b) \text{ } \text{powi} - \text{ int } k) \text{ } A \rangle$

$[\text{THEN } \text{lipschitz-with-onD2}, \text{ OF } \langle x \in A \rangle \langle y \in A \rangle]$

have $\langle \text{dist } (f \ x) \ (f \ y) \leq \text{restriction-}\delta \text{ } \text{TYPE}'(b) \text{ } \text{powi} - \text{ int } k * \text{dist } x \ y \rangle$.

hence $\langle \text{dist } (f \ x) \ (f \ y) * \text{restriction-}\delta \text{ } \text{TYPE}'(b) \wedge k \leq \text{dist } x \ y \rangle$

by (*subst (asm) mult.commute*)

$(\text{drule mult-imp-div-pos-le[rotated]}; \text{simp add: power-int-minus-divide})$

hence $\langle \text{restriction-}\sigma \text{ } \text{TYPE}'(b) \text{ } (\text{Sup } (\text{restriction-related-set } (f \ x) \ (f \ y))) * \text{restriction-}\delta \text{ } \text{TYPE}'(b) \wedge k$

$\leq \text{restriction-}\sigma \text{ } \text{TYPE}'(b) \text{ } (\text{Sup } (\text{restriction-related-set } x \ y)) \rangle$

by (*simp add: dist-restriction-is-bis* $\langle x \neq y \rangle \langle f \ x \neq f \ y \rangle$)

$(\text{drule order-trans[OF - le-restriction-}\sigma], \text{simp})$

from $\text{order-trans[OF restriction-}\sigma\text{-le-restriction-}\sigma\text{-times-pow-restriction-}\delta$

$[\text{of } \langle \text{Sup } (\text{restriction-related-set } (f \ x) \ (f \ y)) \rangle k] \text{ this}]$

have $\langle \text{restriction-}\sigma \text{ } \text{TYPE}'(b) \text{ } (\text{Sup } (\text{restriction-related-set } (f \ x) \ (f \ y)) + k)$

$\leq \text{restriction-}\sigma \text{ } \text{TYPE}'(b) \text{ } (\text{Sup } (\text{restriction-related-set } x \ y)) \rangle$.

hence $\langle \text{Sup } (\text{restriction-related-set } x \ y) \leq \text{Sup } (\text{restriction-related-set } (f \ x) \ (f \ y)) + k \rangle$

using *strict-decseq-def-ter strict-decseq-restriction-}\sigma* **by** *blast*

from $\text{order-trans[OF cSup-upper this]} \text{ have } \langle n \leq \text{Sup } (\text{restriction-related-set } (f \ x) \ (f \ y)) + k \rangle$

by (*simp add:* $\langle x \downarrow n = y \downarrow n \rangle \langle x \neq y \rangle$ *bdd-above-restriction-related-set-iff*)

hence $\langle \text{nat } (\text{int } n + - \text{ int } k) \leq \text{Sup } (\text{restriction-related-set } (f \ x) \ (f \ y)) \rangle$ **by** *linarith*

thus $\langle f \ x \downarrow \text{nat } (\text{int } n + - \text{ int } k) = f \ y \downarrow \text{nat } (\text{int } n + - \text{ int } k) \rangle$

by (*metis not-eq-imp-dist-restriction-is-restriction-}\sigma\text{-Sup-restriction-eq-simplified(2)*)

qed

corollary *lipschitz-with-imp-restriction-shift-neg* :

$\langle \llbracket \text{lipschitz-with } f \text{ } (\text{restriction-}\delta \text{ } \text{TYPE}'(b) \text{ } \text{powi} - \text{ int } k);$

$\bigwedge n. \text{restriction-}\sigma \text{ } \text{TYPE}'(a) \text{ } n \leq \text{restriction-}\sigma \text{ } \text{TYPE}'(b) \text{ } n \rrbracket$

$\implies \text{restriction-shift } f \text{ } (- \text{ int } k) \rangle$

for $f :: \langle 'a :: \text{decseq-restriction-space} \Rightarrow 'b :: \text{at-least-geometric-restriction-space} \rangle$

using *lipschitz-with-on-imp-restriction-shift-neg-on restriction-shift-def*

by *blast*

We obtained that *restriction-shift* implies *lipschitz-with* when $0 \leq k$ and that *lipschitz-with* implies *restriction-shift* when $k \leq 0$. To cover the remaining cases, we give move from *at-least-geometric-restriction-space*

to *geometric-restriction-space*.

theorem *lipschitz-with-on-iff-restriction-shift-on* :

$\langle \text{lipschitz-with-on } f \text{ (restriction-}\delta \text{ TYPE('b) powi } k) \ A \longleftrightarrow \text{restriction-shift-on } f \ k \ A \rangle$

if *same-restriction- σ* : $\langle \text{restriction-}\sigma \ \text{TYPE('b)} = \text{restriction-}\sigma \ \text{TYPE('a)} \rangle$

for $f :: \langle 'a :: \text{decseq-restriction-space} \Rightarrow 'b :: \text{geometric-restriction-space} \rangle$

proof (*rule iffI*)

— We could do a case on k , but both cases are actually handled by the proof required after applying $\llbracket \text{lipschitz-with-on } ?f \ (\delta_{\downarrow} \ \text{TYPE}(?'b) \ \text{powi} \ - \ \text{int } ?k) \ ?A; \bigwedge n. \ \sigma_{\downarrow} \ \text{TYPE}(?'a) \ n \leq \sigma_{\downarrow} \ \text{TYPE}(?'b) \ n \rrbracket \Longrightarrow \text{restriction-shift-on } ?f \ (- \ \text{int } ?k) \ ?A$.

show $\langle \text{restriction-shift-on } f \ k \ A \rangle$ **if** $\langle \text{lipschitz-with-on } f \text{ (restriction-}\delta \ \text{TYPE('b) powi } k) \ A \rangle$

proof (*rule restriction-shift-onI*)

fix $x \ y \ n$ **assume** $\langle x \in A \rangle \langle y \in A \rangle \langle f \ x \neq f \ y \rangle \langle x \downarrow n = y \downarrow n \rangle$

from $\langle f \ x \neq f \ y \rangle$ **have** $\langle x \neq y \rangle$ **by** *blast*

from $\langle \text{lipschitz-with-on } f \text{ (restriction-}\delta \ \text{TYPE('b) powi } k) \ A \rangle$

[*THEN lipschitz-with-onD2, OF* $\langle x \in A \rangle \langle y \in A \rangle$]

have $\langle \text{dist } (f \ x) \ (f \ y) \leq \text{restriction-}\delta \ \text{TYPE('b) powi } k * \text{dist } x \ y \rangle$.

also have $\langle \text{dist } (f \ x) \ (f \ y) = \text{restriction-}\delta \ \text{TYPE('b)} \wedge \text{Sup} \ (\text{restriction-related-set } (f \ x) \ (f \ y)) \rangle$

by (*simp add: dist-restriction-is-bis* $\langle f \ x \neq f \ y \rangle$ *restriction- σ -is*)

also have $\langle \text{dist } x \ y = \text{restriction-}\delta \ \text{TYPE('b)} \wedge \text{Sup} \ (\text{restriction-related-set } x \ y) \rangle$

by (*simp add: dist-restriction-is-bis* $\langle x \neq y \rangle$ *restriction- σ -is same-restriction- σ [symmetric]*)

finally have $\langle \text{restriction-}\delta \ \text{TYPE('b)} \wedge \text{Sup} \ (\text{restriction-related-set } (f \ x) \ (f \ y)) \leq \text{restriction-}\delta \ \text{TYPE('b) powi } k * \text{restriction-}\delta \ \text{TYPE('b)} \wedge \text{Sup} \ (\text{restriction-related-set } x \ y) \rangle$.

also have $\langle \dots = \text{restriction-}\delta \ \text{TYPE('b) powi } (k + \text{Sup} \ (\text{restriction-related-set } x \ y)) \rangle$

by (*subst power-int-add, metis order-less-irrefl zero-less-restriction- δ*) *simp*

finally have $\langle \text{restriction-}\delta \ \text{TYPE('b)} \wedge \text{Sup} \ (\text{restriction-related-set } (f \ x) \ (f \ y)) \leq \dots \rangle$.

moreover have $\langle \text{restriction-}\delta \ \text{TYPE('b) powi } m \leq \text{restriction-}\delta \ \text{TYPE('b) powi } m' \Longrightarrow m' \leq m \rangle$ **for** $m \ m'$

by (*rule ccontr, unfold not-le, drule power-int-strict-decreasing[OF - zero-less-restriction- δ restriction- δ -less-one]*)

(*fold not-le, blast*)

ultimately have $\langle k + \text{Sup} \ (\text{restriction-related-set } x \ y) \leq \text{Sup} \ (\text{restriction-related-set } (f \ x) \ (f \ y)) \rangle$ **by** *simp*

hence $\langle \text{Sup} \ (\text{restriction-related-set } x \ y) \leq \text{Sup} \ (\text{restriction-related-set } (f \ x) \ (f \ y)) - k \rangle$ **by** *simp*

with $\langle x \downarrow n = y \downarrow n \rangle \langle x \neq y \rangle$ *linorder-not-le*

not-eq-imp-dist-restriction-is-restriction- σ -Sup-restriction-eq-simplified(3)

have $\langle n \leq \text{Sup} \ (\text{restriction-related-set } (f \ x) \ (f \ y)) - k \rangle$ **by** *fastforce*

hence $\langle \text{nat } (\text{int } n + k) \leq \text{Sup } (\text{restriction-related-set } (f x) (f y)) \rangle$
by *simp*
thus $\langle f x \downarrow \text{nat } (\text{int } n + k) = f y \downarrow \text{nat } (\text{int } n + k) \rangle$
by (*metis not-eq-imp-dist-restriction-is-restriction- σ -Sup-restriction-eq-simplified(2)*)
qed
next
show $\langle \text{lipschitz-with-on } f \text{ (restriction- δ TYPE('b) powi } k) A \rangle$ **if**
 $\langle \text{restriction-shift-on } f k A \rangle$
proof (*cases k*)
show $\langle k = \text{int } k' \implies \text{lipschitz-with-on } f \text{ (restriction- δ TYPE('b) powi } k) A \rangle$ **for** k'
by (*simp, rule restriction-shift-nonneg-on-imp-lipschitz-with-on*)
(use $\langle \text{restriction-shift-on } f k A \rangle$ same-restriction- σ in simp-all)
next
fix k' **assume** $\langle k = - \text{int } (\text{Suc } k') \rangle$
show $\langle \text{lipschitz-with-on } f \text{ (restriction- δ TYPE('b) powi } k) A \rangle$
proof (*rule lipschitz-with-onI*)
show $\langle 0 \leq \text{restriction- δ TYPE('b) powi } k \rangle$ **by** *simp*
next
fix $x y$ **assume** $\langle x \in A \rangle \langle y \in A \rangle \langle x \neq y \rangle \langle f x \neq f y \rangle$
have $\langle i \in \text{restriction-related-set } x y \implies i - \text{Suc } k' \in \text{restriction-related-set } (f x) (f y) \rangle$ **for** i
using $\langle \text{restriction-shift-on } f k A \rangle$ [*THEN restriction-shift-onD, OF $\langle x \in A \rangle \langle y \in A \rangle$, of i*]
by (*unfold $\langle k = - \text{int } (\text{Suc } k') \rangle$, clarify*) (*metis diff-conv-add-uminus nat-minus-as-int*)
hence $\langle \text{Sup } (\text{restriction-related-set } x y) - \text{Suc } k' \in \text{restriction-related-set } (f x) (f y) \rangle$
using $\langle x \neq y \rangle$ *not-eq-imp-dist-restriction-is-restriction- σ -Sup-restriction-eq-simplified*
by *blast*
hence $\langle \text{Sup } (\text{restriction-related-set } x y) - \text{Suc } k' \leq \text{Sup } (\text{restriction-related-set } (f x) (f y)) \rangle$
by (*simp add: $\langle f x \neq f y \rangle$ bdd-above-restriction-related-set-iff cSup-upper*)
hence $\ast : \langle \text{Sup } (\text{restriction-related-set } x y) + k \leq \text{Sup } (\text{restriction-related-set } (f x) (f y)) \rangle$
by (*simp add: $\langle k = - \text{int } (\text{Suc } k') \rangle$*)
have $\langle \text{dist } (f x) (f y) = \text{restriction- δ TYPE('b) powi } \text{Sup } (\text{restriction-related-set } (f x) (f y)) \rangle$
by (*simp add: $\langle f x \neq f y \rangle$ dist-restriction-is-bis-simplified restriction- σ -is*)
also have $\langle \dots \leq \text{restriction- δ TYPE('b) powi } (\text{Sup } (\text{restriction-related-set } x y) + k) \rangle$
by (*rule power-int-decreasing[OF \ast]; simp*)
(metis order-less-irrefl zero-less-restriction- δ)
also have $\langle \dots = \text{restriction- δ TYPE('b) powi } k \ast \text{restriction- δ TYPE('b) powi } (\text{Sup } (\text{restriction-related-set } x y)) \rangle$
by (*subst power-int-add, metis order-less-irrefl zero-less-restriction- δ*)
simp

also have $\langle \text{restriction-}\delta \text{ TYPE('b) powi (Sup (restriction-related-set } x \ y)) = \text{dist } x \ y \rangle$
by (*simp add: $\langle x \neq y \rangle \text{ dist-restriction-is-bis-simplified}$*
restriction- σ -is-same-restriction- σ [symmetric])
finally show $\langle \text{dist } (f \ x) \ (f \ y) \leq \text{restriction-}\delta \ \text{TYPE('b) powi } k * \text{dist } x \ y \rangle$.
qed
qed
qed

corollary *lipschitz-with-iff-restriction-shift* :
 $\langle \text{restriction-}\sigma \ \text{TYPE('b) = restriction-}\sigma \ \text{TYPE('a)} \implies$
 $\text{lipschitz-with } f \ (\text{restriction-}\delta \ \text{TYPE('b) powi } k) \longleftrightarrow \text{restriction-shift}$
 $f \ k \rangle$
for $f :: \langle 'a :: \text{decseq-restriction-space} \Rightarrow 'b :: \text{geometric-restriction-space} \rangle$
by (*simp add: lipschitz-with-on-iff-restriction-shift-on restriction-shift-def*)

4 Functions

4.1 Restriction Space

instantiation $\langle \text{fun} \rangle :: (\text{type}, \text{restriction-}\sigma) \text{ restriction-}\sigma$
begin

definition *restriction- σ -fun* :: $\langle ('a \Rightarrow 'b) \text{ itself} \Rightarrow \text{nat} \Rightarrow \text{real} \rangle$
where $\langle \text{restriction-}\sigma\text{-fun} - \equiv \text{restriction-}\sigma \ \text{TYPE('b)} \rangle$

instance by *intro-classes*

end

instantiation $\langle \text{fun} \rangle :: (\text{type}, \text{non-decseq-restriction-space}) \text{ non-decseq-restriction-space}$
begin

definition *dist-fun* :: $\langle ['a \Rightarrow 'b, 'a \Rightarrow 'b] \Rightarrow \text{real} \rangle$
where $\langle \text{dist-fun } f \ g \equiv \text{INF } n \in \text{restriction-related-set } f \ g. \text{restriction-}\sigma$
 $\text{TYPE('a} \Rightarrow 'b) \ n \rangle$

definition *uniformity-fun* :: $\langle (('a \Rightarrow 'b) \times ('a \Rightarrow 'b)) \text{ filter} \rangle$
where $\langle \text{uniformity-fun} \equiv \text{INF } e \in \{0 < ..\}. \text{principal } \{(x, y). \text{dist } x \ y$
 $< e\} \rangle$

definition *open-fun* :: $\langle ('a \Rightarrow 'b) \text{ set} \Rightarrow \text{bool} \rangle$
where $\langle \text{open-fun } U \equiv \forall x \in U. \text{eventually } (\lambda(x', y). x' = x \longrightarrow y \in$
 $U) \text{ uniformity} \rangle$

```

instance by intro-classes
  (simp-all add: restriction- $\sigma$ -fun-def dist-fun-def open-fun-def
   uniformity-fun-def restriction- $\sigma$ -tendsto-zero)

end

instance  $\langle \text{fun} \rangle :: (\text{type}, \text{decseq-restriction-space}) \text{decseq-restriction-space}$ 
  by intro-classes (simp add: restriction- $\sigma$ -fun-def decseq-restriction- $\sigma$ )

instance  $\langle \text{fun} \rangle :: (\text{type}, \text{strict-decseq-restriction-space}) \text{strict-decseq-restriction-space}$ 
  by intro-classes
  (simp add: restriction- $\sigma$ -fun-def strict-decseq-restriction- $\sigma$ )

instantiation  $\langle \text{fun} \rangle :: (\text{type}, \text{restriction-}\delta) \text{restriction-}\delta$ 
begin

definition restriction- $\delta$ -fun ::  $\langle ('a \Rightarrow 'b) \text{itself} \Rightarrow \text{real} \rangle$ 
  where  $\langle \text{restriction-}\delta\text{-fun} - \equiv \text{restriction-}\delta \text{TYPE}('b) \rangle$ 

instance by intro-classes (simp-all add: restriction- $\delta$ -fun-def)

end

instance  $\langle \text{fun} \rangle :: (\text{type}, \text{at-least-geometric-restriction-space}) \text{at-least-geometric-restriction-space}$ 
proof intro-classes
  show  $\langle 0 < \text{restriction-}\sigma \text{TYPE}('a \Rightarrow 'b) n \rangle$  for  $n$ 
    by (simp add: restriction- $\sigma$ -fun-def)
  next
    show  $\langle \text{restriction-}\sigma \text{TYPE}('a \Rightarrow 'b) (\text{Suc } n) \leq \text{restriction-}\delta \text{TYPE}('a \Rightarrow 'b) * \text{restriction-}\sigma \text{TYPE}('a \Rightarrow 'b) n \rangle$ 
    for  $n$ 
    by (simp add: restriction- $\sigma$ -le restriction- $\sigma$ -fun-def restriction- $\delta$ -fun-def)
  next
    show  $\langle \text{dist } f g = \text{Inf } (\text{restriction-}\sigma\text{-related-set } f g) \rangle$  for  $f g :: \langle 'a \Rightarrow 'b \rangle$ 
    by (simp add: dist-fun-def)
  qed

instance  $\langle \text{fun} \rangle :: (\text{type}, \text{geometric-restriction-space}) \text{geometric-restriction-space}$ 
proof intro-classes
  show  $\langle \text{restriction-}\sigma \text{TYPE}('a \Rightarrow 'b) n = \text{restriction-}\delta \text{TYPE}('a \Rightarrow 'b) \wedge n \rangle$  for  $n$ 
    by (simp add: restriction- $\sigma$ -fun-def restriction- $\sigma$ -is restriction- $\delta$ -fun-def)
  next
    show  $\langle \text{dist } f g = \text{Inf } (\text{restriction-}\sigma\text{-related-set } f g) \rangle$  for  $f g :: \langle 'a \Rightarrow 'b \rangle$ 

```

by (simp add: dist-fun-def)
qed

lemma *dist-image-le-dist-fun* : $\langle \text{dist } (f \ x) \ (g \ x) \leq \text{dist } f \ g \rangle$
for $f \ g :: \langle 'a \Rightarrow 'b :: \text{non-decseq-restriction-space} \rangle$
proof (unfold dist-restriction-is, rule cInf-superset-mono)
show $\langle \text{restriction-}\sigma\text{-related-set } f \ g \neq \{\} \rangle$ **by** simp
next
show $\langle \text{bdd-below } (\text{restriction-}\sigma\text{-related-set } (f \ x) \ (g \ x)) \rangle$
by (simp add: bounded-imp-bdd-below bounded-restriction- σ -related-set)
next
show $\langle \text{restriction-}\sigma\text{-related-set } f \ g \subseteq \text{restriction-}\sigma\text{-related-set } (f \ x) \ (g \ x) \rangle$
unfolding restriction-fun-def restriction- σ -fun-def
by (simp add: image-def subset-iff) metis
qed

lemma *Sup-dist-image-le-dist-fun* : $\langle (\text{SUP } x. \text{dist } (f \ x) \ (g \ x)) \leq \text{dist } f \ g \rangle$
for $f \ g :: \langle 'a \Rightarrow 'b :: \text{non-decseq-restriction-space} \rangle$
by (simp add: dist-image-le-dist-fun cSUP-least)
— The other inequality will require the additional assumption *decseq*.

context fixes $f \ g :: \langle 'a \Rightarrow 'b :: \text{decseq-restriction-space} \rangle$ **begin**

lemma *reached-dist-fun* : $\langle \exists x. \text{dist } f \ g = \text{dist } (f \ x) \ (g \ x) \rangle$
proof (cases $\langle f = g \rangle$)
show $\langle f = g \implies \exists x. \text{dist } f \ g = \text{dist } (f \ x) \ (g \ x) \rangle$ **by** simp
next
assume $\langle f \neq g \rangle$
let $?n = \langle \text{Sup } (\text{restriction-related-set } f \ g) \rangle$
from not-eq-imp-dist-restriction-is-restriction- σ -Sup-restriction-eq-simplified(2,
3)[OF $\langle f \neq g \rangle$]
obtain x **where** $\langle \forall m \leq ?n. f \ x \downarrow m = g \ x \downarrow m \rangle \langle f \ x \downarrow \text{Suc } ?n \neq g \ x \downarrow \text{Suc } ?n \rangle$
unfolding restriction-fun-def **by** (meson lessI)
hence $\langle \text{dist } (f \ x) \ (g \ x) = \text{restriction-}\sigma \ \text{TYPE}('b) \ ?n \rangle$
by (metis (no-types, lifting) le-neq-implies-less not-less-eq-eq
not-eq-imp-dist-restriction-is-restriction- σ -Sup-restriction-eq-simplified)
with not-eq-imp-dist-restriction-is-restriction- σ -Sup-restriction-eq-simplified(1)
[OF $\langle f \neq g \rangle$, unfolded restriction- σ -fun-def]
have $\langle \text{dist } f \ g = \text{dist } (f \ x) \ (g \ x) \rangle$ **by** simp
thus $\langle \exists x. \text{dist } f \ g = \text{dist } (f \ x) \ (g \ x) \rangle$..
qed

```

lemma dist-fun-eq-Sup-dist-image :  $\langle \text{dist } f \ g = (\text{SUP } x. \text{dist } (f \ x) \ (g \ x)) \rangle$ 
proof (rule order-antisym)
  show  $\langle (\text{SUP } x. \text{dist } (f \ x) \ (g \ x)) \leq \text{dist } f \ g \rangle$  by (fact Sup-dist-image-le-dist-fun)
next
  from reached-dist-fun obtain  $x$  where  $\langle \text{dist } f \ g = \text{dist } (f \ x) \ (g \ x) \rangle$ 
..
thus  $\langle \text{dist } f \ g \leq (\text{SUP } x. \text{dist } (f \ x) \ (g \ x)) \rangle$ 
proof (rule ord-eq-le-trans)
  show  $\langle \text{dist } (f \ x) \ (g \ x) \leq (\text{SUP } x. \text{dist } (f \ x) \ (g \ x)) \rangle$ 
  proof (rule cSup-upper)
    show  $\langle \text{dist } (f \ x) \ (g \ x) \in \text{range } (\lambda x. \text{dist } (f \ x) \ (g \ x)) \rangle$  by simp
  next
    show  $\langle \text{bdd-above } (\text{range } (\lambda x. \text{dist } (f \ x) \ (g \ x))) \rangle$ 
    by (rule bdd-aboveI[of -  $\langle \text{dist } f \ g \rangle$ ]) (auto intro: dist-image-le-dist-fun)
  qed
qed
qed

```

```

lemma fun-restriction-space-Sup-properties :
   $\langle \text{dist } (f \ x) \ (g \ x) \leq \text{dist } f \ g \rangle$ 
   $\langle (\bigwedge x. \text{dist } (f \ x) \ (g \ x) \leq b) \implies \text{dist } f \ g \leq b \rangle$ 
  by (use reached-dist-fun in  $\langle \text{auto simp add: dist-image-le-dist-fun} \rangle$ )

```

end

4.2 Completeness

Actually we can obtain even better: when the instance $'b$ of *decseq-restriction-space* is also an instance of *complete-space*, the type $'a \Rightarrow 'b$ is an instance of *complete-space*.

This is because when $'b$ is an instance of *decseq-restriction-space* (and not only *non-decseq-restriction-space*) the distance between two functions is reached (see $\exists x. \text{dist } ?f \ ?g = \text{dist } (?f \ x) \ (?g \ x)$).

The only remaining thing is to prove that completeness is preserved on higher-order.

```

instance  $\langle \text{fun} \rangle$  :: (type, complete-decseq-restriction-space) complete-decseq-restriction-space
  by intro-classes (fact restriction-chain-imp-restriction-convergent)

```

```

instance  $\langle \text{fun} \rangle$  :: (type, complete-strict-decseq-restriction-space) complete-strict-decseq-restriction-space
  by intro-classes (fact restriction-chain-imp-restriction-convergent)

```

instance $\langle \text{fun} \rangle :: (\text{type}, \text{complete-at-least-geometric-restriction-space})$
complete-at-least-geometric-restriction-space
by *intro-classes (fact restriction-chain-imp-restriction-convergent)*

instance $\langle \text{fun} \rangle :: (\text{type}, \text{complete-geometric-restriction-space})$ *complete-geometric-restriction-space*
by *intro-classes (fact restriction-chain-imp-restriction-convergent)*

4.3 Kind of Extensionality

context **fixes** $f :: \langle 'a :: \text{metric-space}, 'b :: \text{type} \rangle \Rightarrow$
 $\langle 'c :: \text{decseq-restriction-space} \rangle$ **begin**

lemma *lipschitz-with-simplification:*

$\langle \text{lipschitz-with } f \ \alpha \longleftrightarrow (\forall y. \text{lipschitz-with } (\lambda x. f \ x \ y) \ \alpha) \rangle$

proof (*intro iffI allI*)

fix y **assume** $\text{asm} : \langle \text{lipschitz-with } f \ \alpha \rangle$

show $\langle \text{lipschitz-with } (\lambda x. f \ x \ y) \ \alpha \rangle$

proof (*rule lipschitz-withI*)

from $\text{asm}[\text{THEN lipschitz-withD1}]$ **show** $\langle 0 \leq \alpha \rangle$.

next

show $\langle \text{dist } (f \ x1 \ y) \ (f \ x2 \ y) \leq \alpha * \text{dist } x1 \ x2 \rangle$ **for** $x1 \ x2$

by (*rule order-trans[OF - asm[THEN lipschitz-withD2]]*)

(*simp add: dist-image-le-dist-fun*)

qed

next

assume $\text{asm} : \langle \forall y. \text{lipschitz-with } (\lambda x. f \ x \ y) \ \alpha \rangle$

show $\langle \text{lipschitz-with } f \ \alpha \rangle$

proof (*rule lipschitz-withI*)

from $\text{asm}[\text{rule-format}, \text{THEN lipschitz-withD1}]$ **show** $\langle 0 \leq \alpha \rangle$.

next

fix $x1 \ x2$

obtain y **where** $\langle \text{dist } (f \ x1) \ (f \ x2) = \text{dist } (f \ x1 \ y) \ (f \ x2 \ y) \rangle$

by (*meson reached-dist-fun*)

also have $\langle \dots \leq \alpha * \text{dist } x1 \ x2 \rangle$ **by** (*rule asm[rule-format, THEN lipschitz-withD2]*)

finally show $\langle \text{dist } (f \ x1) \ (f \ x2) \leq \alpha * \text{dist } x1 \ x2 \rangle$.

qed

qed

lemma *non-expanding-simplification :*

$\langle \text{non-expanding } f \longleftrightarrow (\forall y. \text{non-expanding } (\lambda x. f \ x \ y)) \rangle$

by (*metis lipschitz-with-simplification non-expanding-on-def*)

lemma *contraction-with-simplification:*

$\langle \text{contraction-with } f \ \alpha \longleftrightarrow (\forall y. \text{contraction-with } (\lambda x. f \ x \ y) \ \alpha) \rangle$

by (*metis contraction-with-on-def lipschitz-with-simplification*)

end

5 Product

The product type $'a \times 'b$ of to metric spaces is already instantiated as a metric space by setting $dist\ x\ y = sqrt\ ((dist\ (fst\ x)\ (fst\ y))^2 + (dist\ (snd\ x)\ (snd\ y))^2)$. Unfortunately, this definition is not compatible with the distance required by the *non-decseq-restriction-space..* We first have to define a new product type with a trivial **typedef**.

5.1 Isomorphic Product Construction

5.1.1 Definition and First Properties

typedef $('a, 'b)\ prod_{max}\ (\langle (- \times_{max} / -) \rangle [21, 20]\ 20) = \langle UNIV :: ('a \times 'b)\ set \rangle$

morphisms *from-prod_{max} to-prod_{max}* **by** *simp*

— Simplifications because the **typedef** is trivial.

declare *from-prod_{max}-inject* [*simp*]
from-prod_{max}-inverse [*simp*]

lemmas *to-prod_{max}-inject-simplified* [*simp*] = *to-prod_{max}-inject* [*simplified*]
and *to-prod_{max}-inverse-simplified*[*simp*] = *to-prod_{max}-inverse*[*simplified*]

lemmas *to-prod_{max}-induct-simplified* = *to-prod_{max}-induct*[*simplified*]
and *to-prod_{max}-cases-simplified* = *to-prod_{max}-cases* [*simplified*]
and *from-prod_{max}-induct-simplified* = *from-prod_{max}-induct*[*simplified*]
and *from-prod_{max}-cases-simplified* = *from-prod_{max}-cases* [*simplified*]

setup-lifting *type-definition-prod_{max}*

lift-definition *Pair_{max}* :: $\langle 'a \Rightarrow 'b \Rightarrow 'a \times_{max} 'b \rangle$ **is** *Pair* .

free-constructors *case-prod_{max}* **for** *Pair_{max} fst_{max} snd_{max}*
by (*metis Pair_{max}.abs-eq from-prod_{max}-inverse surjective-pairing*)
(*metis Pair_{max}.rep-eq prod.inject*)

lemma *fst_{max}-def* : $\langle fst_{max} \equiv map_fun\ from_prod_{max}\ id\ fst \rangle$
by (*intro eq-reflection ext, simp add: fst_{max}-def*,

metis Pair_{max}.rep-eq from-prod_{max}-inverse fst_{max}-def prod.collapse prod_{max}.sel(1)

lemma *fst_{max}-rep-eq* : $\langle \text{fst}_{\text{max}} x = \text{fst} (\text{from-prod}_{\text{max}} x) \rangle$
by (*metis Pair_{max}.rep-eq fst-conv prod_{max}.collapse*)

lemma *fst_{max}-abs-eq [simp]* : $\langle \text{fst}_{\text{max}} (\text{to-prod}_{\text{max}} y) = \text{fst} y \rangle$
by (*metis Pair_{max}.abs-eq prod.exhaust-sel prod_{max}.sel(1)*)

lemma *fst_{max}-transfer [transfer-rule]*: $\langle \text{rel-fun} (\text{pcr-prod}_{\text{max}} (=) (=)) (=) \text{fst} \text{fst}_{\text{max}} \rangle$
by (*metis (mono-tags) Pair_{max}.rep-eq cr-prod_{max}-def fst-conv prod_{max}.collapse prod_{max}.pcr-cr-eq rel-funI*)

lemma *snd_{max}-def* : $\langle \text{snd}_{\text{max}} \equiv \text{map-fun from-prod}_{\text{max}} \text{id} \text{snd} \rangle$
by (*intro eq-reflection ext, simp add: snd_{max}-def, metis Pair_{max}.rep-eq from-prod_{max}-inverse prod.collapse prod_{max}.case*)

lemma *snd_{max}-rep-eq* : $\langle \text{snd}_{\text{max}} x = \text{snd} (\text{from-prod}_{\text{max}} x) \rangle$
by (*metis Pair_{max}.rep-eq prod_{max}.collapse snd-conv*)

lemma *snd_{max}-abs-eq [simp]* : $\langle \text{snd}_{\text{max}} (\text{to-prod}_{\text{max}} y) = \text{snd} y \rangle$
by (*metis Pair_{max}.abs-eq prod.exhaust-sel prod_{max}.sel(2)*)

lemma *snd_{max}-transfer [transfer-rule]* : $\langle \text{rel-fun} (\text{pcr-prod}_{\text{max}} (=) (=)) (=) \text{snd} \text{snd}_{\text{max}} \rangle$
by (*metis (mono-tags, lifting) Pair_{max}.rep-eq cr-prod_{max}-def prod_{max}.collapse prod_{max}.pcr-cr-eq rel-funI snd-conv*)

lemma *case-prod_{max}-def* : $\langle \text{case-prod}_{\text{max}} \equiv \text{map-fun id} (\text{map-fun from-prod}_{\text{max}} \text{id}) \text{case-prod} \rangle$
by (*intro eq-reflection ext, simp add: prod_{max}.case-eq-if fst_{max}-rep-eq snd_{max}-rep-eq split-beta*)

lemma *case-prod_{max}-rep-eq* : $\langle \text{case-prod}_{\text{max}} f p = (\text{case from-prod}_{\text{max}} p \text{ of } (x, y) \Rightarrow f x y) \rangle$
by (*simp add: fst_{max}-rep-eq prod_{max}.case-eq-if snd_{max}-rep-eq split-beta*)

lemma *case-prod_{max}-abs-eq [simp]* : $\langle \text{case-prod}_{\text{max}} f (\text{to-prod}_{\text{max}} q) = (\text{case } q \text{ of } (x, y) \Rightarrow f x y) \rangle$
by (*simp add: prod_{max}.case-eq-if split-beta*)

lemma *case-prod_{max}-transfer [transfer-rule]* : $\langle \text{rel-fun} (=) (\text{rel-fun} (\text{pcr-prod}_{\text{max}} (=) (=)) (=)) \text{case-prod} \text{case-prod}_{\text{max}} \rangle$
by (*simp add: cr-prod_{max}-def fst_{max}-rep-eq prod_{max}.case-eq-if prod_{max}.pcr-cr-eq rel-fun-def snd_{max}-rep-eq split-beta*)

5.2 Syntactic Sugar

The following syntactic sugar is of course recovered from the theory *HOL.Product-Type*.

nonterminal *tuple-args_{max}* and *patterns_{max}*

syntax

$-tuple_{max} \quad :: 'a \Rightarrow tuple_args_{max} \Rightarrow 'a \times_{max} 'b \quad (\langle (1\langle -, / - \rangle) \rangle)$
 $-tuple_arg_{max} \quad :: 'a \Rightarrow tuple_args_{max} \quad (\langle - \rangle)$
 $-tuple_args_{max} \quad :: 'a \Rightarrow tuple_args_{max} \Rightarrow tuple_args_{max} \quad (\langle -, / - \rangle)$
 $-pattern_{max} \quad :: pttrn \Rightarrow patterns_{max} \Rightarrow pttrn \quad (\langle \langle -, / - \rangle \rangle)$
 $\quad \quad \quad :: pttrn \Rightarrow patterns_{max} \quad (\langle - \rangle)$
 $-patterns_{max} \quad :: pttrn \Rightarrow patterns_{max} \Rightarrow patterns_{max} \quad (\langle -, / - \rangle)$

translations

$\langle x, y \rangle \equiv CONST\ Pair_{max}\ x\ y$
 $-pattern_{max}\ x\ y \equiv CONST\ Pair_{max}\ x\ y$
 $-patterns_{max}\ x\ y \equiv CONST\ Pair_{max}\ x\ y$
 $-tuple_{max}\ x\ (-tuple_args_{max}\ y\ z) \equiv -tuple_{max}\ x\ (-tuple_arg_{max}\ (-tuple_{max}\ y\ z))$
 $\lambda \langle x, y, zs \rangle. b \equiv CONST\ case_prod_{max}\ (\lambda x\ \langle y, zs \rangle. b)$
 $\lambda \langle x, y \rangle. b \equiv CONST\ case_prod_{max}\ (\lambda x\ y. b)$
 $-abs\ (CONST\ Pair_{max}\ x\ y)\ t \rightarrow \lambda \langle x, y \rangle. t$
 — This rule accommodates tuples in *case C ...* $\langle x, y \rangle \dots \Rightarrow \dots$:
 The $\langle x, y \rangle$ is parsed as *Pair_{max} x y* because it is *logic*, not *pttrn*.

With this syntactic sugar, one can write *case a of* $\langle b, c, d, e \rangle \Rightarrow \langle c, d \rangle, \lambda \langle y, u \rangle. a, \lambda \langle a, b \rangle. \langle a, b, c, d, e \rangle, \lambda \langle a, b, c \rangle. a, \dots$ as for the type $'a \times 'b$.

lemmas *to-prod_{max}-tuple* [simp] = *Pair_{max}.abs-eq*[*symmetric*]
and *from-prod_{max}-tuple_{max}* [simp] = *Pair_{max}.rep-eq*

5.3 Product

We first redo the work of *Restriction-Spaces.Restricton-Spaces-Prod*.

instantiation *prod_{max}* :: (*restriction, restriction*) *restriction*
begin

lift-definition *restriction-prod_{max}* :: $\langle 'a \times_{max} 'b \Rightarrow nat \Rightarrow 'a \times_{max} 'b \rangle$ is $\langle (\downarrow) \rangle$.

lemma *restriction-prod_{max}-def'* : $\langle p \downarrow n = \langle fst_{max}\ p \downarrow n, snd_{max}\ p \downarrow n \rangle \rangle$
by *transfer* (*simp add: restriction-prod-def*)

instance by (*intro-classes, transfer, simp*)

end

instance $prod_{max} :: (restriction\text{-}space, restriction\text{-}space) restriction\text{-}space$
by $(intro\text{-}classes; transfer)$ $(simp\text{-}all\ add: ex\text{-}not\text{-}restriction\text{-}related)$

instantiation $prod_{max} :: (restriction\text{-}\sigma, restriction\text{-}\sigma) restriction\text{-}\sigma$
begin

definition $restriction\text{-}\sigma\text{-}prod_{max} :: \langle 'a \times_{max} 'b \rangle itself \Rightarrow nat \Rightarrow real \rangle$
where $\langle restriction\text{-}\sigma\text{-}prod_{max} - n \equiv$
 $max (restriction\text{-}\sigma\ TYPE('a) n) (restriction\text{-}\sigma\ TYPE('b) n) \rangle$

instance by $intro\text{-}classes$
end

instantiation $prod_{max} :: (non\text{-}decseq\text{-}restriction\text{-}space, non\text{-}decseq\text{-}restriction\text{-}space)$
 $non\text{-}decseq\text{-}restriction\text{-}space$
begin

definition $dist\text{-}prod_{max} :: \langle ['a \times_{max} 'b, 'a \times_{max} 'b] \Rightarrow real \rangle$
where $\langle dist\text{-}prod_{max} f g \equiv INF n \in restriction\text{-}related\text{-}set f g. re\text{-}$
 $striction\text{-}\sigma\ TYPE('a \times_{max} 'b) n \rangle$

definition $uniformity\text{-}prod_{max} :: \langle (('a \times_{max} 'b) \times 'a \times_{max} 'b) filter \rangle$
where $\langle uniformity\text{-}prod_{max} \equiv INF e \in \{0 < ..\}. principal \{(x, y). dist$
 $x y < e\} \rangle$

definition $open\text{-}prod_{max} :: \langle ('a \times_{max} 'b) set \Rightarrow bool \rangle$
where $\langle open\text{-}prod_{max} U \equiv \forall x \in U. eventually (\lambda(x', y). x' = x \longrightarrow$
 $y \in U) uniformity \rangle$

instance

proof $intro\text{-}classes$

show $\langle restriction\text{-}\sigma\ TYPE('a \times_{max} 'b) \longrightarrow 0 \rangle$
by $(rule\ real\text{-}tendsto\text{-}sandwich$
 $[of \langle \lambda n. 0 \rangle - - \langle \lambda n. restriction\text{-}\sigma\ TYPE('a) n + restriction\text{-}\sigma$
 $TYPE('b) n \rangle])$
 $(simp\text{-}all\ add: order\text{-}less\text{-}imp\text{-}le\ restriction\text{-}\sigma\text{-}prod_{max}\text{-}def\ max\text{-}def$
 $restriction\text{-}\sigma\text{-}tendsto\text{-}zero\ tendsto\text{-}add\text{-}zero)$
qed $(simp\text{-}all\ add: uniformity\text{-}prod_{max}\text{-}def\ open\text{-}prod_{max}\text{-}def$
 $restriction\text{-}\sigma\text{-}prod_{max}\text{-}def\ max\text{-}def\ dist\text{-}prod_{max}\text{-}def)$

end

instance $prod_{max} :: (decseq\text{-}restriction\text{-}space, decseq\text{-}restriction\text{-}space)$
 $decseq\text{-}restriction\text{-}space$
proof $intro\text{-}classes$

```

show ⟨decseq (restriction-σ TYPE('a ×max 'b))⟩
proof (intro decseq-SucI)
  show ⟨restriction-σ TYPE('a ×max 'b) (Suc n) ≤ restriction-σ
TYPE('a ×max 'b) n⟩ for n
  using decseq-SucD[of ⟨restriction-σ TYPE('a)⟩ n]
  decseq-SucD[of ⟨restriction-σ TYPE('b)⟩ n]
  by (auto simp add: restriction-σ-prodmax-def decseq-restriction-σ)
qed
qed

```

```

instance prodmax :: (strict-decseq-restriction-space, strict-decseq-restriction-space)
strict-decseq-restriction-space
proof intro-classes
  show ⟨strict-decseq (restriction-σ TYPE('a ×max 'b))⟩
  proof (intro strict-decseq-SucI)
    show ⟨restriction-σ TYPE('a ×max 'b) (Suc n) < restriction-σ
TYPE('a ×max 'b) n⟩ for n
    using strict-decseq-SucD[of ⟨restriction-σ TYPE('a)⟩ n]
    strict-decseq-SucD[of ⟨restriction-σ TYPE('b)⟩ n]
    by (auto simp add: restriction-σ-prodmax-def strict-decseq-restriction-σ)
  qed
qed

```

```

instantiation prodmax :: (restriction-δ, restriction-δ) restriction-δ
begin

```

```

definition restriction-δ-prodmax :: ⟨('a ×max 'b) itself ⇒ real⟩
  where ⟨restriction-δ-prodmax - ≡ max (restriction-δ TYPE('a))
(restriction-δ TYPE('b))⟩

```

```

instance by intro-classes (simp-all add: restriction-δ-prodmax-def max-def)

```

```

end

```

```

instance prodmax :: (at-least-geometric-restriction-space, at-least-geometric-restriction-space)
at-least-geometric-restriction-space
proof intro-classes
  show ⟨0 < restriction-σ TYPE('a ×max 'b) n⟩ for n by simp
next
  show ⟨restriction-σ TYPE('a ×max 'b) (Suc n)
≤ restriction-δ TYPE('a ×max 'b) * restriction-σ TYPE('a
×max 'b) n⟩ for n
  by (auto intro: order-trans[OF restriction-σ-le]
simp add: restriction-δ-prodmax-def mult-mono' restriction-σ-prodmax-def)
next
  show ⟨dist p1 p2 = Inf (restriction-σ-related-set p1 p2)⟩ for p1 p2
:: ⟨'a ×max 'b⟩
  by (simp add: dist-prodmax-def)

```

qed

```
instance prodmax :: (geometric-restriction-space, geometric-restriction-space)
geometric-restriction-space
proof intro-classes
  show ⟨restriction-σ TYPE('a ×max 'b) n = restriction-δ TYPE('a
×max 'b) ^ n⟩ for n
  by (simp add: restriction-σ-prodmax-def restriction-σ-is restric-
tion-δ-prodmax-def max-def)
(meson nle-le power-mono zero-le-restriction-δ)
next
  show ⟨dist p1 p2 = Inf (restriction-σ-related-set p1 p2)⟩ for p1 p2
:: ⟨'a ×max 'b⟩
  by (simp add: dist-prodmax-def)
qed
```

```
lemma max-dist-projections-le-dist-prodmax :
⟨max (dist (fstmax p1) (fstmax p2)) (dist (sndmax p1) (sndmax p2))
≤ dist p1 p2⟩
proof (unfold dist-restriction-is max-def, split if-split, intro conjI impI)
  show ⟨Inf (restriction-σ-related-set (sndmax p1) (sndmax p2)) ≤ Inf
(restriction-σ-related-set p1 p2)⟩
  proof (rule cINF-superset-mono[OF nonempty-restriction-related-set])
    show ⟨bdd-below (restriction-σ-related-set (sndmax p1) (sndmax
p2))⟩
    by (meson bdd-belowI2 zero-le-restriction-σ)
  qed (simp-all add: subset-iff add: restriction-prodmax-def' restric-
tion-σ-prodmax-def)
next
  show ⟨Inf (restriction-σ-related-set (fstmax p1) (fstmax p2)) ≤ Inf
(restriction-σ-related-set p1 p2)⟩
  proof (rule cINF-superset-mono[OF nonempty-restriction-related-set])
    show ⟨bdd-below (restriction-σ-related-set (fstmax p1) (fstmax p2))⟩
    by (meson bdd-belowI2 zero-le-restriction-σ)
  qed (simp-all add: subset-iff add: restriction-prodmax-def' restric-
tion-σ-prodmax-def)
qed
```

5.4 Completeness

5.4.1 Preliminaries

default-sort non-decseq-restriction-space — Otherwise we should al-
ways specify.

```
lemma restriction-σ-prodmax-commute :
⟨restriction-σ TYPE('b ×max 'a) = restriction-σ TYPE('a ×max
```

'b)⟩

unfolding *restriction- σ -prod_{max}-def* **by** (*rule ext*) *simp*

definition *dist-left-prod_{max}* :: $\langle ('a \times_{max} 'b) \text{ itself} \Rightarrow 'a \Rightarrow 'a \Rightarrow \text{real} \rangle$
where $\langle \text{dist-left-prod}_{max} - x y \equiv \text{INF } n \in \text{restriction-related-set } x$
 $y. \text{restriction-}\sigma \text{ TYPE}('a \times_{max} 'b) n \rangle$

definition *dist-right-prod_{max}* :: $\langle ('a \times_{max} 'b) \text{ itself} \Rightarrow 'b \Rightarrow 'b \Rightarrow \text{real} \rangle$

where $\langle \text{dist-right-prod}_{max} - x y \equiv \text{INF } n \in \text{restriction-related-set } x$
 $y. \text{restriction-}\sigma \text{ TYPE}('a \times_{max} 'b) n \rangle$

lemma *dist-right-prod_{max}-is-dist-left-prod_{max}* :
 $\langle \text{dist-right-prod}_{max} \text{ TYPE}('b \times_{max} 'a) = \text{dist-left-prod}_{max} \text{ TYPE}('a$
 $\times_{max} 'b) \rangle$

unfolding *dist-left-prod_{max}-def* *dist-right-prod_{max}-def*

by (*subst restriction- σ -prod_{max}-commute*) *simp*

lemma *dist-le-dist-left-prod_{max}* : $\langle \text{dist } x y \leq \text{dist-left-prod}_{max} \text{ TYPE}('a$
 $\times_{max} 'b) x y \rangle$

proof (*unfold dist-left-prod_{max}-def* *dist-restriction-is*,

rule cINF-mono[OF nonempty-restriction-related-set[of x y]])

show $\langle \text{bdd-below } (\text{restriction-}\sigma\text{-related-set } x y) \rangle$

by (*meson bdd-belowI2 zero-le-restriction- σ*)

next

show $\langle m \in \text{restriction-related-set } x y \Longrightarrow$

$\exists n \in \text{restriction-related-set } x y. \sigma_{\downarrow} \text{ TYPE}('a) n \leq \sigma_{\downarrow} \text{ TYPE}('a$

$\times_{max} 'b) m \rangle$ **for** *m*

by (*metis max.cobounded1 restriction- σ -prod_{max}-def*)

qed

lemma *dist-le-dist-right-prod_{max}* : $\langle \text{dist } x y \leq \text{dist-right-prod}_{max} \text{ TYPE}('b$
 $\times_{max} 'a) x y \rangle$

by (*simp add: dist-le-dist-left-prod_{max} dist-right-prod_{max}-is-dist-left-prod_{max}*)

lemma

fixes *p1 p2* :: $\langle 'a :: \text{decseq-restriction-space} \times_{max} 'b :: \text{decseq-restriction-space} \rangle$

shows *dist-prod_{max}-le-max-dist-left-prod_{max}-dist-right-prod_{max}* :

$\langle \text{dist } p1 p2 \leq \text{max } (\text{dist-left-prod}_{max} \text{ TYPE}('a \times_{max} 'b) (\text{fst}_{max}$
 $p1) (\text{fst}_{max} p2))$

$(\text{dist-right-prod}_{max} \text{ TYPE}('a \times_{max} 'b) (\text{snd}_{max}$
 $p1) (\text{snd}_{max} p2)) \rangle$

proof –

interpret *left* : *DecseqRestrictionSpace* $\langle (\downarrow) \rangle \langle (=) \rangle \langle \text{UNIV} \rangle$

$\langle \text{restriction-}\sigma \text{ TYPE}('a \times_{max} 'b) \rangle \langle \text{dist-left-prod}_{max} \text{ TYPE}('a$
 $\times_{max} 'b) \rangle$

by *unfold-locales*
 (*simp-all add: restriction- σ -tendsto-zero dist-left-prod_{max}-def*
decseq-restriction- σ)

interpret *right* : *DecseqRestrictionSpace* $\langle (\downarrow) \rangle \langle (=) \rangle \langle UNIV :: 'b \text{ set} \rangle$
 $\langle \text{restriction-}\sigma \text{ TYPE}('a \times_{max} 'b) \rangle \langle \text{dist-right-prod}_{max} \text{ TYPE}('a$
 $\times_{max} 'b) \rangle$

by *unfold-locales*
 (*simp-all add: restriction- σ -tendsto-zero dist-right-prod_{max}-def*
decseq-restriction- σ)

show $\langle \text{dist } p1 \ p2 \leq \text{max } (\text{dist-left-prod}_{max} \text{ TYPE}('a \times_{max} 'b)$
 $(fst_{max} \ p1) \ (fst_{max} \ p2)) \rangle$
 $(\text{dist-right-prod}_{max} \ \text{TYPE}('a \times_{max} 'b) \ (snd_{max}$
 $p1) \ (snd_{max} \ p2)) \rangle$

by (*auto simp add: dist-restriction-is-bis left.dist-restriction-is-bis*
right.dist-restriction-is-bis prod_{max}.expand restriction-prod_{max}-def)
 (*smt (verit, best) Collect-cong nle-le restriction-related-le*)

qed

default-sort *type* — Back to normal.

5.4.2 Complete Restriction Space

When the instances *'a* and *'b* of *decseq-restriction-space* are also instances of *complete-space*, the type *'a* \times_{max} *'b* is an instance of *complete-space*.

lemma *restriction-chain-prod_{max}-iff* :
 $\langle \text{restriction-chain } \sigma \longleftrightarrow \text{restriction-chain } (\lambda n. \text{fst}_{max} (\sigma \ n)) \wedge$
 $\text{restriction-chain } (\lambda n. \text{snd}_{max} (\sigma \ n)) \rangle$

by (*simp add: restriction-chain-def, transfer*)
 (*metis fst-conv prod.collapse restriction-prod-def snd-conv*)

lemma *restriction-tendsto-prod_{max}-iff* :
 $\langle \sigma \ -\downarrow \rightarrow \Sigma \longleftrightarrow (\lambda n. \text{fst}_{max} (\sigma \ n)) \ -\downarrow \rightarrow \text{fst}_{max} \ \Sigma \wedge (\lambda n. \text{snd}_{max}$
 $(\sigma \ n)) \ -\downarrow \rightarrow \text{snd}_{max} \ \Sigma \rangle$

by (*simp add: restriction-tendsto-def, transfer, simp add: restric-*
tion-prod-def)
 (*meson nle-le order.trans*)

lemma *restriction-convergent-prod_{max}-iff* :
 $\langle \text{restriction-convergent } \sigma \longleftrightarrow \text{restriction-convergent } (\lambda n. \text{fst}_{max} (\sigma$
 $n)) \wedge$
 $\text{restriction-convergent } (\lambda n. \text{snd}_{max} (\sigma \ n)) \rangle$

by (*simp add: restriction-convergent-def restriction-tendsto-prod_{max}-iff*)
 (*metis prod_{max}.sel*)

```

instance prodmax :: (complete-decseq-restriction-space, complete-decseq-restriction-space)
  complete-decseq-restriction-space
proof (intro-classes, transfer)
  fix σ :: ⟨nat ⇒ 'a ×max 'b⟩ assume ⟨chain↓ σ⟩
  hence ⟨chain↓ (λn. fstmax (σ n))⟩ ⟨chain↓ (λn. sndmax (σ n))⟩
    by (simp-all add: restriction-chain-prodmax-iff)
  hence ⟨convergent↓ (λn. fstmax (σ n))⟩ ⟨convergent↓ (λn. sndmax
    (σ n))⟩
    by (simp-all add: restriction-chain-imp-restriction-convergent)
  thus ⟨convergent↓ σ⟩
    by (simp add: restriction-convergent-prodmax-iff)
qed

```

```

instance prodmax :: (complete-strict-decseq-restriction-space, complete-strict-decseq-restriction-space)
  complete-strict-decseq-restriction-space
by intro-classes (simp add: restriction-chain-imp-restriction-convergent)

```

```

instance prodmax :: (complete-at-least-geometric-restriction-space, complete-at-least-geometric-restriction-space)
  complete-at-least-geometric-restriction-space
by intro-classes (simp add: restriction-chain-imp-restriction-convergent)

```

```

instance prodmax :: (complete-geometric-restriction-space, complete-geometric-restriction-space)
  complete-geometric-restriction-space
by intro-classes (simp add: restriction-chain-imp-restriction-convergent)

```

When the types 'a and 'b share the same restriction sequence, we have the following equality.

```

lemma same-restriction-σ-imp-restriction-σ-prodmax-is [simp] :
  ⟨restriction-σ TYPE('b :: non-decseq-restriction-space) =
  restriction-σ TYPE('a :: non-decseq-restriction-space) ⇒
  restriction-σ TYPE('a ×max 'b) = restriction-σ TYPE('a)⟩
unfolding restriction-σ-prodmax-def by simp

```

```

lemma same-restriction-σ-imp-dist-prodmax-eq-max-dist-projections :
  ⟨dist p1 p2 = max (dist (fstmax p1) (fstmax p2)) (dist (sndmax p1)
  (sndmax p2))⟩
if same-restriction-σ [simp] : ⟨restriction-σ TYPE('b) = restriction-σ
  TYPE('a)⟩
for p1 p2 :: ⟨'a :: decseq-restriction-space ×max 'b :: decseq-restriction-space⟩
proof (rule order-antisym)
  have ⟨dist-left-prodmax TYPE('a ×max 'b) (fstmax p1) (fstmax p2)
  = dist (fstmax p1) (fstmax p2)⟩
    by (simp add: dist-left-prodmax-def dist-restriction-is)
  moreover have ⟨dist-right-prodmax TYPE('a ×max 'b) (sndmax
  p1) (sndmax p2) = dist (sndmax p1) (sndmax p2)⟩

```

```

    by (simp add: dist-right-prodmax-def dist-restriction-is)
  ultimately show ⟨dist p1 p2 ≤ max (dist (fstmax p1) (fstmax p2))
    (dist (sndmax p1) (sndmax p2))⟩
  by (metis dist-prodmax-le-max-dist-left-prodmax-dist-right-prodmax)
next
  show ⟨max (dist (fstmax p1) (fstmax p2)) (dist (sndmax p1) (sndmax
    p2)) ≤ dist p1 p2⟩
  by (fact max-dist-projections-le-dist-prodmax)
qed

```

Now, one can write things like $v \langle x, y \rangle. f \langle x, y \rangle$.

We could of course imagine more support for $'a \times_{max} 'b$ type, but as restriction spaces are intended to be used without recourse to metric spaces, we have not undertaken this task for the time being.

6 Main entry Point