

Relative Security

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Abstract

This entry formalizes the notion of relative security, which can be used to model transient execution vulnerabilities in the style of Spectre and Meltdown. The notion was introduced in the CSF 2023 paper “Relative Security: Formally Modeling and (Dis)Proving Resilience Against Semantic Optimization Vulnerabilities” by Brijesh Dongol, Matt Griffin, Andrei Popescu and Jamie Wright [1].

It defines two versions of relative security: a finitary one (restricted to finite traces), and an infinitary one (working with both finite and infinite traces). It formalizes unwinding methods for verifying relative security in both the finitary and infinitary versions, and proves their soundness. The proof of soundness in the infinitary case is a substantial application of Isabelle’s corecursion and coinduction infrastructure.

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1 Finitary Relative Security

This theory formalizes the finitary version of relative security, more precisely the notion expressed in terms of finite traces.

```
theory Relative-Security-fin
imports Preliminaries/Transition-System
begin
```

```
declare Let-def[simp]
```

```
no-notation relcomp (infixr O 75)
no-notation relcompp (infixr OO 75)
```

1.1 Finite-trace versions of leakage models and attacker models

```
locale Leakage-Mod-fin = System-Mod istate validTrans final
for istate :: 'state  $\Rightarrow$  bool and validTrans :: 'state  $\times$  'state  $\Rightarrow$  bool and final ::
'state  $\Rightarrow$  bool
+
fixes S :: 'state list  $\Rightarrow$  'secret list
and A :: 'state trace  $\Rightarrow$  'act list
and O :: 'state trace  $\Rightarrow$  'obs list
and leakVia :: 'state list  $\Rightarrow$  'state list  $\Rightarrow$  'leak  $\Rightarrow$  bool
```

```
locale Attacker-Mod-fin = System-Mod istate validTrans final
for istate :: 'state  $\Rightarrow$  bool and validTrans :: 'state  $\times$  'state  $\Rightarrow$  bool and final ::
'state  $\Rightarrow$  bool
+
fixes S :: 'state list  $\Rightarrow$  'secret list
and A :: 'state trace  $\Rightarrow$  'act list
and O :: 'state trace  $\Rightarrow$  'obs list
begin
```

```
fun leakVia :: 'state list  $\Rightarrow$  'state list  $\Rightarrow$  'secret list  $\times$  'secret list  $\Rightarrow$  bool
where
leakVia tr tr' (sl,sl') = (S tr = sl  $\wedge$  S tr' = sl'  $\wedge$  A tr = A tr'  $\wedge$  O tr  $\neq$  O tr')
```

```
lemmas leakVia-def = leakVia.simps
```

```
end
```

```
sublocale Attacker-Mod-fin < Leakage-Mod-fin
```

where $leakVia = leakVia$
 by *standard*

1.2 Locales for increasingly concrete notions of finitary relative security

locale *Relative-Security''-fin* =
Van: Leakage-Mod-fin *istateV* *validTransV* *finalV* *SV* *AV* *OV* *leakViaV*
 +
Opt: Leakage-Mod-fin *istateO* *validTransO* *finalO* *SO* *AO* *OO* *leakViaO*
for *validTransV* :: 'stateV × 'stateV ⇒ bool
and *istateV* :: 'stateV ⇒ bool **and** *finalV* :: 'stateV ⇒ bool
and *SV* :: 'stateV list ⇒ 'secret list
and *AV* :: 'stateV trace ⇒ 'actV list
and *OV* :: 'stateV trace ⇒ 'obsV list
and *leakViaV* :: 'stateV list ⇒ 'stateV list ⇒ 'leak ⇒ bool

and *validTransO* :: 'stateO × 'stateO ⇒ bool
and *istateO* :: 'stateO ⇒ bool **and** *finalO* :: 'stateO ⇒ bool
and *SO* :: 'stateO list ⇒ 'secret list
and *AO* :: 'stateO trace ⇒ 'actO list
and *OO* :: 'stateO trace ⇒ 'obsO list
and *leakViaO* :: 'stateO list ⇒ 'stateO list ⇒ 'leak ⇒ bool

and *corrState* :: 'stateV ⇒ 'stateO ⇒ bool
begin

definition *rsecure* :: bool **where**

rsecure ≡ ∀ l s1 tr1 s2 tr2.

istateO s1 ∧ *Opt.validFromS* s1 tr1 ∧ *Opt.completedFrom* s1 tr1 ∧
istateO s2 ∧ *Opt.validFromS* s2 tr2 ∧ *Opt.completedFrom* s2 tr2 ∧
leakViaO tr1 tr2 l

→

(∃ sv1 trv1 sv2 trv2.

istateV sv1 ∧ *istateV* sv2 ∧ *corrState* sv1 s1 ∧ *corrState* sv2 s2 ∧
Van.validFromS sv1 trv1 ∧ *Van.completedFrom* sv1 trv1 ∧
Van.validFromS sv2 trv2 ∧ *Van.completedFrom* sv2 trv2 ∧
leakViaV trv1 trv2 l)

end

locale *Relative-Security'-fin* =

Van: Attacker-Mod-fin *istateV* *validTransV* *finalV* *SV* *AV* *OV*

+

Opt: Attacker-Mod-fin *istateO* *validTransO* *finalO* *SO* *AO* *OO*

for *validTransV* :: 'stateV × 'stateV ⇒ bool

and *istateV* :: 'stateV ⇒ bool **and** *finalV* :: 'stateV ⇒ bool

and *SV* :: 'stateV list ⇒ 'secret list

```

and AV :: 'stateV trace  $\Rightarrow$  'actV list
and OV :: 'stateV trace  $\Rightarrow$  'obsV list

and validTransO :: 'stateO  $\times$  'stateO  $\Rightarrow$  bool
and istateO :: 'stateO  $\Rightarrow$  bool and finalO :: 'stateO  $\Rightarrow$  bool
and SO :: 'stateO list  $\Rightarrow$  'secret list
and AO :: 'stateO trace  $\Rightarrow$  'actO list
and OO :: 'stateO trace  $\Rightarrow$  'obsO list
and corrState :: 'stateV  $\Rightarrow$  'stateO  $\Rightarrow$  bool

sublocale Relative-Security'-fin < Relative-Security''-fin
where leakViaV = Van.leakVia and leakViaO = Opt.leakVia
by standard

context Relative-Security'-fin
begin

lemma rsecure-def2:
rsecure  $\longleftrightarrow$ 
( $\forall$  s1 tr1 s2 tr2.
  istateO s1  $\wedge$  Opt.validFromS s1 tr1  $\wedge$  Opt.completedFrom s1 tr1  $\wedge$ 
  istateO s2  $\wedge$  Opt.validFromS s2 tr2  $\wedge$  Opt.completedFrom s2 tr2  $\wedge$ 
  AO tr1 = AO tr2  $\wedge$  OO tr1  $\neq$  OO tr2
 $\longrightarrow$ 
  ( $\exists$  sv1 trv1 sv2 trv2.
    istateV sv1  $\wedge$  istateV sv2  $\wedge$  corrState sv1 s1  $\wedge$  corrState sv2 s2  $\wedge$ 
    Van.validFromS sv1 trv1  $\wedge$  Van.completedFrom sv1 trv1  $\wedge$ 
    Van.validFromS sv2 trv2  $\wedge$  Van.completedFrom sv2 trv2  $\wedge$ 
    SV trv1 = SO tr1  $\wedge$  SV trv2 = SO tr2  $\wedge$ 
    AV trv1 = AV trv2  $\wedge$  OV trv1  $\neq$  OV trv2))

unfolding rsecure-def
unfolding Van.leakVia-def Opt.leakVia-def
by auto metis

end

locale Statewise-Attacker-Mod = System-Mod istate validTrans final
for istate :: 'state  $\Rightarrow$  bool and validTrans :: 'state  $\times$  'state  $\Rightarrow$  bool and final ::
'state  $\Rightarrow$  bool
+
fixes
  isSec :: 'state  $\Rightarrow$  bool and getSec :: 'state  $\Rightarrow$  'secret
and
  isInt :: 'state  $\Rightarrow$  bool and getInt :: 'state  $\Rightarrow$  'act  $\times$  'obs
assumes final-not-isInt:  $\bigwedge s. \text{final } s \implies \neg \text{isInt } s$ 

```

and *final-not-isSec*: $\bigwedge s. \text{final } s \implies \neg \text{isSec } s$
begin

definition *getAct* :: 'state \Rightarrow 'act **where**
getAct = *fst* o *getInt*

definition *getObs* :: 'state \Rightarrow 'obs **where**
getObs = *snd* o *getInt*

definition *eqObs trn1 trn2* \equiv
(*isInt trn1* \longleftrightarrow *isInt trn2*) \wedge (*isInt trn1* \longrightarrow *getObs trn1* = *getObs trn2*)

definition *eqAct trn1 trn2* \equiv
(*isInt trn1* \longleftrightarrow *isInt trn2*) \wedge (*isInt trn1* \longrightarrow *getAct trn1* = *getAct trn2*)

definition *A* :: 'state trace \Rightarrow 'act list **where**
A tr \equiv *filtermap isInt getAct (butlast tr)*

sublocale *A*: *FiltermapBL isInt getAct A*
apply standard unfolding *A-def* ..

definition *O* :: 'state trace \Rightarrow 'obs list **where**
O tr \equiv *filtermap isInt getObs (butlast tr)*

sublocale *O*: *FiltermapBL isInt getObs O*
apply standard unfolding *O-def* ..

definition *S* :: 'state list \Rightarrow 'secret list **where**
S tr \equiv *filtermap isSec getSec (butlast tr)*

sublocale *S*: *FiltermapBL isSec getSec S*
apply standard unfolding *S-def* ..

end

sublocale *Statewise-Attacker-Mod* < *Attacker-Mod-fin*
where *S* = *S* **and** *A* = *A* **and** *O* = *O*
by standard

locale *Rel-Sec* =

```

  Van: Statewise-Attacker-Mod istateV validTransV finalV isSecV getSecV isIntV
  getIntV
+
  Opt: Statewise-Attacker-Mod istateO validTransO finalO isSecO getSecO isIntO
  getIntO
  for validTransV :: 'stateV × 'stateV ⇒ bool
  and istateV :: 'stateV ⇒ bool and finalV :: 'stateV ⇒ bool
  and isSecV :: 'stateV ⇒ bool and getSecV :: 'stateV ⇒ 'secret
  and isIntV :: 'stateV ⇒ bool and getIntV :: 'stateV ⇒ 'actV × 'obsV

  and validTransO :: 'stateO × 'stateO ⇒ bool
  and istateO :: 'stateO ⇒ bool and finalO :: 'stateO ⇒ bool
  and isSecO :: 'stateO ⇒ bool and getSecO :: 'stateO ⇒ 'secret
  and isIntO :: 'stateO ⇒ bool and getIntO :: 'stateO ⇒ 'actO × 'obsO

  and corrState :: 'stateV ⇒ 'stateO ⇒ bool

sublocale Rel-Sec < Relative-Security'-fin
where SV = Van.S and AV = Van.A and OV = Van.O
and SO = Opt.S and AO = Opt.A and OO = Opt.O
by standard

context Rel-Sec
begin

abbreviation getObsV :: 'stateV ⇒ 'obsV where getObsV ≡ Van.getObs
abbreviation getActV :: 'stateV ⇒ 'actV where getActV ≡ Van.getAct
abbreviation getObsO :: 'stateO ⇒ 'obsO where getObsO ≡ Opt.getObs
abbreviation getActO :: 'stateO ⇒ 'actO where getActO ≡ Opt.getAct

abbreviation reachV where reachV ≡ Van.reach
abbreviation reachO where reachO ≡ Opt.reach

abbreviation completedFromV :: 'stateV ⇒ 'stateV list ⇒ bool where complet-
  edFromV ≡ Van.completedFrom
abbreviation completedFromO :: 'stateO ⇒ 'stateO list ⇒ bool where complet-
  edFromO ≡ Opt.completedFrom

lemmas completedFromV-def = Van.completedFrom-def
lemmas completedFromO-def = Opt.completedFrom-def

lemma rsecure-def3:
rsecure ⟷
(∀ s1 tr1 s2 tr2.
  istateO s1 ∧ Opt.validFromS s1 tr1 ∧ completedFromO s1 tr1 ∧
  istateO s2 ∧ Opt.validFromS s2 tr2 ∧ completedFromO s2 tr2 ∧
  Opt.A tr1 = Opt.A tr2 ∧ Opt.O tr1 ≠ Opt.O tr2 ∧

```

$(isIntO\ s1 \wedge isIntO\ s2 \longrightarrow getActO\ s1 = getActO\ s2)$
 \longrightarrow
 $(\exists sv1\ trv1\ sv2\ trv2.$
 $\quad istateV\ sv1 \wedge istateV\ sv2 \wedge corrState\ sv1\ s1 \wedge corrState\ sv2\ s2 \wedge$
 $\quad Van.validFromS\ sv1\ trv1 \wedge completedFromV\ sv1\ trv1 \wedge$
 $\quad Van.validFromS\ sv2\ trv2 \wedge completedFromV\ sv2\ trv2 \wedge$
 $\quad Van.S\ trv1 = Opt.S\ tr1 \wedge Van.S\ trv2 = Opt.S\ tr2 \wedge$
 $\quad Van.A\ trv1 = Van.A\ trv2 \wedge Van.O\ trv1 \neq Van.O\ trv2))$
unfolding *rsecure-def2* **apply** (*intro iff-allI iffI impI*)
subgoal by auto
subgoal
by *clarsimp (metis (full-types) Opt.A.Cons-unfold*
 $Opt.completed-Cons\ Opt.final-not-isInt$
 $Simple-Transition-System.validFromS-Cons-iff$
 $completedFromO-def\ list.sel(1)\ neq-Nil-conv)$.

definition $eqSec\ trnO\ trnA \equiv$
 $(isSecV\ trnO = isSecO\ trnA) \wedge (isSecV\ trnO \longrightarrow getSecV\ trnO = getSecO\ trnA)$

lemma *eqSec-S-Cons'*:
 $eqSec\ trnO\ trnA \implies$
 $(Van.S\ (trnO \# trO') = Opt.S\ (trnA \# trA')) \implies Van.S\ trO' = Opt.S\ trA'$
apply(*cases trO' = []*)
subgoal apply(*cases trA' = []*)
subgoal by auto
subgoal unfolding *eqSec-def* **by auto** .
subgoal apply(*cases trA' = []*)
subgoal by auto
subgoal unfolding *eqSec-def* **by auto** . .

lemma *eqSec-S-Cons[simp]*:
 $eqSec\ trnO\ trnA \implies trO' = [] \longleftrightarrow trA' = [] \implies$
 $(Van.S\ (trnO \# trO') = Opt.S\ (trnA \# trA')) \longleftrightarrow (Van.S\ trO' = Opt.S\ trA')$
apply(*cases trO' = []*)
subgoal apply(*cases trA' = []*)
subgoal by auto
subgoal unfolding *eqSec-def* **by auto** .
subgoal apply(*cases trA' = []*)
subgoal by auto
subgoal unfolding *eqSec-def* **by auto** . .

end

```

locale Relative-Security-Determ =
  Rel-Sec
    validTransV istateV finalV isSecV getSecV isIntV getIntV
    validTransO istateO finalO isSecO getSecO isIntO getIntO
    corrState
  +
  System-Mod-Deterministic istateV validTransV finalV nextO
    for validTransV :: 'stateV × 'stateV ⇒ bool
    and istateV :: 'stateV ⇒ bool
    and finalV :: 'stateV ⇒ bool
    and nextO :: 'stateV ⇒ 'stateV
    and isSecV :: 'stateV ⇒ bool and getSecV :: 'stateV ⇒ 'secret
    and isIntV :: 'stateV ⇒ bool and getIntV :: 'stateV ⇒ 'actV × 'obsV
    and validTransO :: 'stateO × 'stateO ⇒ bool
    and istateO :: 'stateO ⇒ bool
    and finalO :: 'stateO ⇒ bool
    and isSecO :: 'stateO ⇒ bool and getSecO :: 'stateO ⇒ 'secret
    and isIntO :: 'stateO ⇒ bool and getIntO :: 'stateO ⇒ 'actO × 'obsO
    and corrState :: 'stateV ⇒ 'stateO ⇒ bool

end

```

2 Relative Security

This theory formalizes the general notion of relative security, applicable to possibly infinite traces.

```

theory Relative-Security
imports Relative-Security-fin Preliminaries/Trivial
begin

```

```

no-notation relcomp (infixr O 75)
no-notation relcompp (infixr OO 75)

```

2.1 Leakage models and attacker models

```

locale Leakage-Mod = System-Mod istate validTrans final
for istate :: 'state ⇒ bool and validTrans :: 'state × 'state ⇒ bool and final ::
  'state ⇒ bool
  +
fixes S :: 'state llist ⇒ 'secret llist
and A :: 'state ltrace ⇒ 'act llist
and O :: 'state ltrace ⇒ 'obs llist
and leakVia :: 'state llist ⇒ 'state llist ⇒ 'leak ⇒ bool

```

```

locale Attacker-Mod = System-Mod istate validTrans final

```



```

for istate :: 'state  $\Rightarrow$  bool and validTrans :: 'state  $\times$  'state  $\Rightarrow$  bool and final ::
'state  $\Rightarrow$  bool
+
fixes S :: 'state llist  $\Rightarrow$  'secret llist
and A :: 'state ltrace  $\Rightarrow$  'act llist
and O :: 'state ltrace  $\Rightarrow$  'obs llist
begin

fun leakVia :: 'state llist  $\Rightarrow$  'state llist  $\Rightarrow$  'secret llist  $\times$  'secret llist  $\Rightarrow$  bool
where
leakVia tr tr' (sl, sl') = (S tr = sl  $\wedge$  S tr' = sl'  $\wedge$  A tr = A tr'  $\wedge$  O tr  $\neq$  O tr')

lemmas leakVia-def = leakVia.simps

end

sublocale Attacker-Mod < Leakage-Mod
where leakVia = leakVia
by standard

```

2.2 Locales for increasingly concrete notions of relative security

```

locale Relative-Security'' =
  Van: Leakage-Mod istateV validTransV finalV SV AV OV leakViaV
+
  Opt: Leakage-Mod istateO validTransO finalO SO AO OO leakViaO
for validTransV :: 'stateV  $\times$  'stateV  $\Rightarrow$  bool
and istateV :: 'stateV  $\Rightarrow$  bool and finalV :: 'stateV  $\Rightarrow$  bool
and SV :: 'stateV llist  $\Rightarrow$  'secret llist
and AV :: 'stateV ltrace  $\Rightarrow$  'actV llist
and OV :: 'stateV ltrace  $\Rightarrow$  'obsV llist
and leakViaV :: 'stateV llist  $\Rightarrow$  'stateV llist  $\Rightarrow$  'leak  $\Rightarrow$  bool

and validTransO :: 'stateO  $\times$  'stateO  $\Rightarrow$  bool
and istateO :: 'stateO  $\Rightarrow$  bool and finalO :: 'stateO  $\Rightarrow$  bool
and SO :: 'stateO llist  $\Rightarrow$  'secret llist
and AO :: 'stateO ltrace  $\Rightarrow$  'actO llist
and OO :: 'stateO ltrace  $\Rightarrow$  'obsO llist
and leakViaO :: 'stateO llist  $\Rightarrow$  'stateO llist  $\Rightarrow$  'leak  $\Rightarrow$  bool

and corrState :: 'stateV  $\Rightarrow$  'stateO  $\Rightarrow$  bool
begin

```

definition *lrsecure* :: bool **where**

```

lrsecure  $\equiv$   $\forall$  l s1 tr1 s2 tr2.
  istateO s1  $\wedge$  Opt.linvalidFromS s1 tr1  $\wedge$  Opt.lcompletedFrom s1 tr1  $\wedge$ 
  istateO s2  $\wedge$  Opt.linvalidFromS s2 tr2  $\wedge$  Opt.lcompletedFrom s2 tr2  $\wedge$ 

```

$leakViaO\ tr1\ tr2\ l$
 \longrightarrow
 $(\exists\ sv1\ trv1\ sv2\ trv2.$
 $\quad istateV\ sv1 \wedge istateV\ sv2 \wedge corrState\ sv1\ s1 \wedge corrState\ sv2\ s2 \wedge$
 $\quad Van.linvalidFromS\ sv1\ trv1 \wedge Van.lcompletedFrom\ sv1\ trv1 \wedge$
 $\quad Van.linvalidFromS\ sv2\ trv2 \wedge Van.lcompletedFrom\ sv2\ trv2 \wedge$
 $\quad leakViaV\ trv1\ trv2\ l)$

end

locale *Relative-Security'* =
 $Van: Attacker-Mod\ istateV\ validTransV\ finalV\ SV\ AV\ OV$
 $+$
 $Opt: Attacker-Mod\ istateO\ validTransO\ finalO\ SO\ AO\ OO$
for $validTransV :: 'stateV \times 'stateV \Rightarrow bool$
and $istateV :: 'stateV \Rightarrow bool$ **and** $finalV :: 'stateV \Rightarrow bool$
and $SV :: 'stateV\ llist \Rightarrow 'secret\ llist$
and $AV :: 'stateV\ ltrace \Rightarrow 'actV\ llist$
and $OV :: 'stateV\ ltrace \Rightarrow 'obsV\ llist$

and $validTransO :: 'stateO \times 'stateO \Rightarrow bool$
and $istateO :: 'stateO \Rightarrow bool$ **and** $finalO :: 'stateO \Rightarrow bool$
and $SO :: 'stateO\ llist \Rightarrow 'secret\ llist$
and $AO :: 'stateO\ ltrace \Rightarrow 'actO\ llist$
and $OO :: 'stateO\ ltrace \Rightarrow 'obsO\ llist$
and $corrState :: 'stateV \Rightarrow 'stateO \Rightarrow bool$

sublocale *Relative-Security' < Relative-Security''*
where $leakViaV = Van.leakVia$ **and** $leakViaO = Opt.leakVia$
by *standard*

context *Relative-Security'*
begin

lemma *lrsecure-def2:*

$lrsecure \longleftrightarrow$
 $(\forall\ s1\ tr1\ s2\ tr2.$
 $\quad istateO\ s1 \wedge Opt.linvalidFromS\ s1\ tr1 \wedge Opt.lcompletedFrom\ s1\ tr1 \wedge$
 $\quad istateO\ s2 \wedge Opt.linvalidFromS\ s2\ tr2 \wedge Opt.lcompletedFrom\ s2\ tr2 \wedge$
 $\quad AO\ tr1 = AO\ tr2 \wedge OO\ tr1 \neq OO\ tr2$
 \longrightarrow
 $(\exists\ sv1\ trv1\ sv2\ trv2.$
 $\quad istateV\ sv1 \wedge istateV\ sv2 \wedge corrState\ sv1\ s1 \wedge corrState\ sv2\ s2 \wedge$

```

    Van.lvalidFromS sv1 trv1 ∧ Van.lcompletedFrom sv1 trv1 ∧
    Van.lvalidFromS sv2 trv2 ∧ Van.lcompletedFrom sv2 trv2 ∧
    SV trv1 = SO tr1 ∧ SV trv2 = SO tr2 ∧
    AV trv1 = AV trv2 ∧ OV trv1 ≠ OV trv2))
unfolding lrsecure-def
unfolding Van.lleakVia-def Opt.lleakVia-def
by auto metis

end

context Statewise-Attacker-Mod begin

definition lA :: 'state ltrace ⇒ 'act llist where
lA tr ≡ lfiltermap isInt getAct (lbutlast tr)

sublocale lA: LfiltermapBL isInt getAct lA
apply standard unfolding lA-def ..

lemma lA: lcompletedFrom s tr ⇒ lA tr = lmap getAct (lfilter isInt tr)
apply(cases lfinite tr)
  subgoal unfolding lA.lmap-lfilter lbutlast-def
  by simp (metis final-not-isInt lbutlast-lfinite lcompletedFrom-def lfilter-llist-of lfiltermap-lmap-lfilter lfinite-lfiltermap-butlast llast-llist-of llist-of-list-of lmap-llist-of)
  subgoal unfolding lA.lmap-lfilter lbutlast-def by auto .

definition lO :: 'state ltrace ⇒ 'obs llist where
lO tr ≡ lfiltermap isInt getObs (lbutlast tr)

sublocale lO: LfiltermapBL isInt getObs lO
apply standard unfolding lO-def ..

lemma lO: lcompletedFrom s tr ⇒ lO tr = lmap getObs (lfilter isInt tr)
apply(cases lfinite tr)
  subgoal unfolding lO.lmap-lfilter lbutlast-def
  by simp (metis List-Filtermap.filtermap-def butlast.simps(1) filtermap-butlast final-not-isInt lcompletedFrom-def lfilter-llist-of llist-of-list-of lmap-llist-of)
  subgoal unfolding lO.lmap-lfilter lbutlast-def by auto .

definition lS :: 'state llist ⇒ 'secret llist where
lS tr ≡ lfiltermap isSec getSec (lbutlast tr)

```

sublocale *lS*: *LfiltermapBL isSec getSec lS*
apply *standard unfolding lS-def ..*

lemma *lS*: *lcompletedFrom s tr \implies lS tr = lmap getSec (lfilter isSec tr)*
apply(*cases lfinite tr*)
subgoal unfolding *lS.lmap-lfilter lbutlast-def*
by simp (*metis List-Filtermap.filtermap-def filtermap-butlast final-not-isSec lcompletedFrom-def lfilter-llist-of llist-of-eq-LNil-conv llist-of-list-of lmap-llist-of*)
subgoal unfolding *lS.lmap-lfilter lbutlast-def by auto .*

end

sublocale *Statewise-Attacker-Mod < Attacker-Mod*
where *S = lS and A = lA and O = lO*
by standard

sublocale *Rel-Sec < Relative-Security'*
where *SV = Van.lS and AV = Van.lA and OV = Van.lO*
and *SO = Opt.lS and AO = Opt.lA and OO = Opt.lO*
by standard

context *Rel-Sec*
begin

abbreviation *lcompletedFromV :: 'stateV \Rightarrow 'stateV llist \Rightarrow bool* **where** *lcompletedFromV \equiv Van.lcompletedFrom*

abbreviation *lcompletedFromO :: 'stateO \Rightarrow 'stateO llist \Rightarrow bool* **where** *lcompletedFromO \equiv Opt.lcompletedFrom*

lemma *eqSec-lS-Cons'*:
eqSec trnO trnA \implies
(Van.lS (trnO \$ trO') = Opt.lS (trnA \$ trA')) \implies Van.lS trO' = Opt.lS trA'
apply(*cases trO' = []*)
subgoal apply(*cases trA' = []*)
subgoal by auto
subgoal unfolding *eqSec-def by auto .*
subgoal apply(*cases trA' = []*)
subgoal by auto
subgoal unfolding *eqSec-def by auto . .*

lemma *eqSec-lS-Cons[simp]*:
eqSec trnO trnA \implies trO' = [] \longleftrightarrow trA' = [] \implies

```

  (Van.lS (trnO $ trO') = Opt.lS (trnA $ trA'))  $\longleftrightarrow$  (Van.lS trO' = Opt.lS trA')
apply(cases trO' = [])
  subgoal apply(cases trA' = [])
    subgoal by auto
    subgoal unfolding eqSec-def by auto .
    subgoal apply(cases trA' = [])
      subgoal by auto
      subgoal unfolding eqSec-def by auto . .
end
end

```

3 Unwinding Proof Method for Finitary Relative Security

This theory formalizes the notion of unwinding for finitary relative security, and proves its soundness.

```

theory Unwinding-fin
imports Relative-Security
begin

```

3.1 The types and operators underlying unwinding: status, matching operators, etc.

```

context Rel-Sec
begin

```

```

datatype status = Eq | Diff

```

```

fun updStat :: status  $\Rightarrow$  bool  $\times$  'a  $\Rightarrow$  bool  $\times$  'a  $\Rightarrow$  status where
  updStat Eq (True,a) (True,a') = (if a = a' then Eq else Diff)
| updStat stat - - = stat

```

```

definition sstatO' statO sv1 sv2 = updStat statO (isIntV sv1, getObsV sv1)
(isIntV sv2, getObsV sv2)

```

```

definition sstatA' statA s1 s2 = updStat statA (isIntO s1, getObsO s1) (isIntO
s2, getObsO s2)

```

```

definition initCond ::

```

```

(enat  $\Rightarrow$  'stateO  $\Rightarrow$  'stateO  $\Rightarrow$  status  $\Rightarrow$  'stateV  $\Rightarrow$  'stateV  $\Rightarrow$  status  $\Rightarrow$  bool)  $\Rightarrow$ 
bool where

```

```

initCond  $\Delta \equiv \forall s1 s2.$ 

```

```

  istateO s1  $\wedge$  istateO s2

```

```

 $\longrightarrow$ 

```

```

  ( $\exists sv1 sv2.$  istateV sv1  $\wedge$  istateV sv2  $\wedge$  corrState sv1 s1  $\wedge$  corrState sv2 s2)

```

$\wedge \Delta \infty s1\ s2\ Eq\ sv1\ sv2\ Eq)$

definition *match1-1* $\Delta\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO \equiv$
 $\exists sv1'.\ validTransV\ (sv1,sv1') \wedge$
 $\Delta \infty s1'\ s2\ statA\ sv1'\ sv2\ statO$

definition *match1-12* $\Delta\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO \equiv$
 $\exists sv1'\ sv2'.$
let $statO' = sstatO'\ statO\ sv1\ sv2$ *in*
 $validTransV\ (sv1,sv1') \wedge$
 $validTransV\ (sv2,sv2') \wedge$
 $\Delta \infty s1'\ s2\ statA\ sv1'\ sv2'\ statO'$

definition *match1* $\Delta\ s1\ s2\ statA\ sv1\ sv2\ statO \equiv$
 $\neg\ isIntO\ s1 \longrightarrow$
 $(\forall s1'.\ validTransO\ (s1,s1')$
 \longrightarrow
 $(\neg\ isSecO\ s1 \wedge \Delta \infty s1'\ s2\ statA\ sv1\ sv2\ statO) \vee$
 $(eqSec\ sv1\ s1 \wedge \neg\ isIntV\ sv1 \wedge match1-1\ \Delta\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO) \vee$
 $(eqSec\ sv1\ s1 \wedge \neg\ isSecV\ sv2 \wedge Van.eqAct\ sv1\ sv2 \wedge match1-12\ \Delta\ s1\ s1'\ s2$
 $statA\ sv1\ sv2\ statO))$

lemmas *match1-defs* = *match1-def match1-1-def match1-12-def*

lemma *match1-1-mono*:
 $\Delta \leq \Delta' \implies match1-1\ \Delta\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO \implies match1-1\ \Delta'\ s1\ s1'\ s2$
 $statA\ sv1\ sv2\ statO$
unfolding *le-fun-def match1-1-def* **by** *auto*

lemma *match1-12-mono*:
 $\Delta \leq \Delta' \implies match1-12\ \Delta\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO \implies match1-12\ \Delta'\ s1\ s1'$
 $s2\ statA\ sv1\ sv2\ statO$
unfolding *le-fun-def match1-12-def* **by** *fastforce*

lemma *match1-mono*:
assumes $\Delta \leq \Delta'$
shows *match1* $\Delta\ s1\ s2\ statA\ sv1\ sv2\ statO \implies match1\ \Delta'\ s1\ s2\ statA\ sv1\ sv2$
 $statO$
unfolding *match1-def* **apply** *clarify subgoal for s1'* **apply**(*erule alle[of - s1']*)
using *match1-1-mono*[*OF* *assms*, *of* *s1 s1' s2 statA sv1 sv2 statO*]
match1-12-mono[*OF* *assms*, *of* *s1 s1' s2 statA sv1 sv2 statO*]
assms[*unfolded le-fun-def, rule-format, of - s1' s2 statA sv1 sv2 statO*]
by *auto* .

definition *match2-1* $\Delta s1 s2 s2' statA sv1 sv2 statO \equiv$
 $\exists sv2'. \text{validTransV } (sv2, sv2') \wedge$
 $\Delta \infty s1 s2' statA sv1 sv2' statO$

definition *match2-12* $\Delta s1 s2 s2' statA sv1 sv2 statO \equiv$
 $\exists sv1' sv2'.$
 $\text{let } statO' = sstatO' statO sv1 sv2 \text{ in}$
 $\text{validTransV } (sv1, sv1') \wedge$
 $\text{validTransV } (sv2, sv2') \wedge$
 $\Delta \infty s1 s2' statA sv1' sv2' statO'$

definition *match2* $\Delta s1 s2 statA sv1 sv2 statO \equiv$
 $\neg \text{isIntO } s2 \longrightarrow$
 $(\forall s2'. \text{validTransO } (s2, s2')$
 \longrightarrow
 $(\neg \text{isSecO } s2 \wedge \Delta \infty s1 s2' statA sv1 sv2 statO) \vee$
 $(\text{eqSec } sv2 s2 \wedge \neg \text{isIntV } sv2 \wedge \text{match2-1 } \Delta s1 s2 s2' statA sv1 sv2 statO) \vee$
 $(\neg \text{isSecV } sv1 \wedge \text{eqSec } sv2 s2 \wedge \text{Van.eqAct } sv1 sv2 \wedge \text{match2-12 } \Delta s1 s2 s2'$
 $\text{statA sv1 sv2 statO}))$

lemmas *match2-defs* = *match2-def match2-1-def match2-12-def*

lemma *match2-1-mono*:
 $\Delta \leq \Delta' \implies \text{match2-1 } \Delta s1 s1' s2 statA sv1 sv2 statO \implies \text{match2-1 } \Delta' s1 s1' s2$
 $\text{statA sv1 sv2 statO}$
unfolding *le-fun-def match2-1-def* **by** *auto*

lemma *match2-12-mono*:
 $\Delta \leq \Delta' \implies \text{match2-12 } \Delta s1 s1' s2 statA sv1 sv2 statO \implies \text{match2-12 } \Delta' s1 s1'$
 $s2 statA sv1 sv2 statO$
unfolding *le-fun-def match2-12-def* **by** *fastforce*

lemma *match2-mono*:
assumes $\Delta \leq \Delta'$
shows $\text{match2 } \Delta s1 s2 statA sv1 sv2 statO \implies \text{match2 } \Delta' s1 s2 statA sv1 sv2$
 statO
unfolding *match2-def* **apply** *clarify subgoal for s2'* **apply**(*erule alle[of - s2']*)
using *match2-1-mono[OF assms, of s1 s2 s2' statA sv1 sv2 statO]*
 $\text{match2-12-mono[OF assms, of s1 s2 s2' statA sv1 sv2 statO]}$
 $\text{assms[unfolded le-fun-def, rule-format, of - s1 s2' statA sv1 sv2 statO]}$
by *auto* .

definition *match12-1* $\Delta s1' s2' statA' sv1 sv2 statO \equiv$
 $\exists sv1'. \text{validTransV } (sv1, sv1') \wedge$
 $\Delta \infty s1' s2' statA' sv1' sv2 statO$

definition *match12-2* $\Delta s1' s2' statA' sv1 sv2 statO \equiv$

$\exists sv2'. \text{validTransV } (sv2, sv2') \wedge$
 $\Delta \infty s1' s2' \text{statA}' sv1 sv2' \text{statO}$

definition *match12-12* $\Delta s1' s2' \text{statA}' sv1 sv2 \text{statO} \equiv$
 $\exists sv1' sv2'.$

let $\text{statO}' = \text{sstatO}' \text{statO } sv1 sv2$ *in*
 $\text{validTransV } (sv1, sv1') \wedge$
 $\text{validTransV } (sv2, sv2') \wedge$
 $(\text{statA}' = \text{Diff} \longrightarrow \text{statO}' = \text{Diff}) \wedge$
 $\Delta \infty s1' s2' \text{statA}' sv1' sv2' \text{statO}'$

definition *match12* $\Delta s1 s2 \text{statA } sv1 sv2 \text{statO} \equiv$
 $\forall s1' s2'.$

let $\text{statA}' = \text{sstatA}' \text{statA } s1 s2$ *in*
 $\text{validTransO } (s1, s1') \wedge$
 $\text{validTransO } (s2, s2') \wedge$
 $\text{Opt.eqAct } s1 s2 \wedge$
 $\text{isIntO } s1 \wedge \text{isIntO } s2$
 \longrightarrow
 $(\neg \text{isSecO } s1 \wedge \neg \text{isSecO } s2 \wedge (\text{statA} = \text{statA}' \vee \text{statO} = \text{Diff}) \wedge \Delta \infty s1' s2'$
 $\text{statA}' sv1 sv2 \text{statO})$
 \vee
 $(\neg \text{isSecO } s2 \wedge \text{eqSec } sv1 s1 \wedge \neg \text{isIntV } sv1 \wedge$
 $(\text{statA} = \text{statA}' \vee \text{statO} = \text{Diff}) \wedge$
 $\text{match12-1 } \Delta s1' s2' \text{statA}' sv1 sv2 \text{statO})$
 \vee
 $(\neg \text{isSecO } s1 \wedge \text{eqSec } sv2 s2 \wedge \neg \text{isIntV } sv2 \wedge$
 $(\text{statA} = \text{statA}' \vee \text{statO} = \text{Diff}) \wedge$
 $\text{match12-2 } \Delta s1' s2' \text{statA}' sv1 sv2 \text{statO})$
 \vee
 $(\text{eqSec } sv1 s1 \wedge \text{eqSec } sv2 s2 \wedge \text{Van.eqAct } sv1 sv2 \wedge$
 $\text{match12-12 } \Delta s1' s2' \text{statA}' sv1 sv2 \text{statO})$

lemmas *match12-defs* = *match12-def match12-1-def match12-2-def match12-12-def*

lemma *match12-simpleI*:

assumes $\bigwedge s1' s2' \text{statA}'.$

$\text{statA}' = \text{sstatA}' \text{statA } s1 s2 \implies$

$\text{validTransO } (s1, s1') \implies$

$\text{validTransO } (s2, s2') \implies$

$\text{Opt.eqAct } s1 s2 \implies$

$(\neg \text{isSecO } s1 \wedge \neg \text{isSecO } s2 \wedge (\text{statA} = \text{statA}' \vee \text{statO} = \text{Diff}) \wedge \Delta \infty s1' s2'$
 $\text{statA}' sv1 sv2 \text{statO})$

\vee

$(\text{eqSec } sv1 s1 \wedge \text{eqSec } sv2 s2 \wedge \text{Van.eqAct } sv1 sv2 \wedge$

$\text{match12-12 } \Delta s1' s2' \text{statA}' sv1 sv2 \text{statO})$

shows *match12* $\Delta s1 s2 \text{statA } sv1 sv2 \text{statO}$

using *assms unfolding match12-def Let-def by blast*

lemma *match12-1-mono*:

$\Delta \leq \Delta' \implies \text{match12-1 } \Delta \ s1' \ s2' \ \text{statA}' \ \text{sv1} \ \text{sv2} \ \text{statO} \implies \text{match12-1 } \Delta' \ s1' \ s2' \ \text{statA}' \ \text{sv1} \ \text{sv2} \ \text{statO}$

unfolding *le-fun-def match12-1-def* **by** *auto*

lemma *match12-2-mono*:

$\Delta \leq \Delta' \implies \text{match12-2 } \Delta \ s1 \ s2' \ \text{statA}' \ \text{sv1} \ \text{sv2} \ \text{statO} \implies \text{match12-2 } \Delta' \ s1 \ s2' \ \text{statA}' \ \text{sv1} \ \text{sv2} \ \text{statO}$

unfolding *le-fun-def match12-2-def* **by** *auto*

lemma *match12-12-mono*:

$\Delta < \Delta' \implies \text{match12-12 } \Delta \ s1' \ s2' \ \text{statA}' \ \text{sv1} \ \text{sv2} \ \text{statO} \implies \text{match12-12 } \Delta' \ s1' \ s2' \ \text{statA}' \ \text{sv1} \ \text{sv2} \ \text{statO}$

unfolding *le-fun-def match12-12-def* **by** *fastforce*

lemma *match12-mono*:

assumes $\Delta \leq \Delta'$

shows $\text{match12 } \Delta \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO} \implies \text{match12 } \Delta' \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}$

unfolding *match12-def* **apply** *clarify* **subgoal** **for** $s1' \ s2'$ **apply** (*erule* *allE*[*of* - $s1'$]) **apply** (*erule* *allE*[*of* - $s2'$])

using *match12-1-mono*[*OF* *assms*, *of* $s1' \ s2' - \text{sv1} \ \text{sv2} \ \text{statO}$]

match12-2-mono[*OF* *assms*, *of* $s1' \ s2' - \text{sv1} \ \text{sv2} \ \text{statO}$]

match12-12-mono[*OF* *assms*, *of* $s1' \ s2' - \text{sv1} \ \text{sv2} \ \text{statO}$]

assms[*unfolded* *le-fun-def*, *rule-format*, *of* - $s1' \ s2'$

sstatA' statA s1 s2 sv1 sv2 statO]

by *simp metis* .

definition *match* $\Delta \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO} \equiv$

match1 $\Delta \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}$

\wedge

match2 $\Delta \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}$

\wedge

match12 $\Delta \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}$

lemmas *match-defs* = *match1-def* *match2-def* *match12-def*

lemmas *match-deep-defs* = *match1-defs* *match2-defs* *match12-defs*

lemma *match-mono*:

assumes $\Delta \leq \Delta'$

shows $\text{match } \Delta \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO} \implies \text{match } \Delta' \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}$

unfolding *match-def* **using** *match1-mono*[*OF* *assms*] *match2-mono*[*OF* *assms*] *match12-mono*[*OF* *assms*] **by** *auto*

definition *move-1* $\Delta w s1 s2 statA sv1 sv2 statO \equiv$
 $\exists sv1'. \text{validTransV } (sv1, sv1') \wedge$
 $\Delta w s1 s2 statA sv1' sv2 statO$

definition *move-2* $\Delta w s1 s2 statA sv1 sv2 statO \equiv$
 $\exists sv2'. \text{validTransV } (sv2, sv2') \wedge$
 $\Delta w s1 s2 statA sv1 sv2' statO$

definition *move-12* $\Delta w s1 s2 statA sv1 sv2 statO \equiv$
 $\exists sv1' sv2'.$
 $\text{let } statO' = sstatO' statO sv1 sv2 \text{ in}$
 $\text{validTransV } (sv1, sv1') \wedge \text{validTransV } (sv2, sv2') \wedge$
 $\Delta w s1 s2 statA sv1' sv2' statO'$

definition *proact* $\Delta w s1 s2 statA sv1 sv2 statO \equiv$
 $(\neg \text{isSecV } sv1 \wedge \neg \text{isIntV } sv1 \wedge \text{move-1 } \Delta w s1 s2 statA sv1 sv2 statO)$
 \vee
 $(\neg \text{isSecV } sv2 \wedge \neg \text{isIntV } sv2 \wedge \text{move-2 } \Delta w s1 s2 statA sv1 sv2 statO)$
 \vee
 $(\neg \text{isSecV } sv1 \wedge \neg \text{isSecV } sv2 \wedge \text{Van.eqAct } sv1 sv2 \wedge \text{move-12 } \Delta w s1 s2 statA$
 $sv1 sv2 statO)$

lemmas *proact-defs* = *proact-def move-1-def move-2-def move-12-def*

lemma *move-1-mono*:
 $\Delta \leq \Delta' \implies \text{move-1 } \Delta \text{ meas } s1 s2 statA sv1 sv2 statO \implies \text{move-1 } \Delta' \text{ meas } s1 s2$
 $statA sv1 sv2 statO$
unfolding *le-fun-def move-1-def* **by** *auto*

lemma *move-2-mono*:
 $\Delta \leq \Delta' \implies \text{move-2 } \Delta \text{ meas } s1 s2 statA sv1 sv2 statO \implies \text{move-2 } \Delta' \text{ meas } s1 s2$
 $statA sv1 sv2 statO$
unfolding *le-fun-def move-2-def* **by** *auto*

lemma *move-12-mono*:
 $\Delta \leq \Delta' \implies \text{move-12 } \Delta \text{ meas } s1 s2 statA sv1 sv2 statO \implies \text{move-12 } \Delta' \text{ meas } s1$
 $s2 statA sv1 sv2 statO$
unfolding *le-fun-def move-12-def* **by** *fastforce*

lemma *proact-mono*:
assumes $\Delta \leq \Delta'$
shows $\text{proact } \Delta \text{ meas } s1 s2 statA sv1 sv2 statO \implies \text{proact } \Delta' \text{ meas } s1 s2 statA$
 $sv1 sv2 statO$
unfolding *proact-def* **using** *move-1-mono[OF assms] move-2-mono[OF assms]*
move-12-mono[OF assms] **by** *auto*

3.2 The definition of unwinding

definition *unwindCond* ::

$(enat \Rightarrow 'stateO \Rightarrow 'stateO \Rightarrow status \Rightarrow 'stateV \Rightarrow 'stateV \Rightarrow status \Rightarrow bool) \Rightarrow bool$

where

$unwindCond \Delta \equiv \forall w \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO.$
 $reachO \ s1 \wedge reachO \ s2 \wedge reachV \ sv1 \wedge reachV \ sv2 \wedge$
 $\Delta \ w \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO$
 \longrightarrow
 $(finalO \ s1 \longleftrightarrow finalO \ s2) \wedge (finalV \ sv1 \longleftrightarrow finalO \ s1) \wedge (finalV \ sv2 \longleftrightarrow finalO$
 $s2)$
 \wedge
 $(statA = Eq \longrightarrow (isIntO \ s1 \longleftrightarrow isIntO \ s2))$
 \wedge
 $((\exists v < w. \ proact \ \Delta \ v \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO)$
 \vee
 $match \ \Delta \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO$
 $)$

lemma *unwindCond-simpleI*:

assumes

$\bigwedge w \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO.$
 $reachO \ s1 \Longrightarrow reachO \ s2 \Longrightarrow reachV \ sv1 \Longrightarrow reachV \ sv2 \Longrightarrow$
 $\Delta \ w \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO$
 \Longrightarrow
 $(finalO \ s1 \longleftrightarrow finalO \ s2) \wedge (finalV \ sv1 \longleftrightarrow finalO \ s1) \wedge (finalV \ sv2 \longleftrightarrow finalO$
 $s2)$

and

$\bigwedge w \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO.$
 $reachO \ s1 \Longrightarrow reachO \ s2 \Longrightarrow reachV \ sv1 \Longrightarrow reachV \ sv2 \Longrightarrow$
 $\Delta \ w \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO \Longrightarrow statA = Eq$
 \Longrightarrow
 $isIntO \ s1 \longleftrightarrow isIntO \ s2$

and

$\bigwedge w \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO.$
 $reachO \ s1 \Longrightarrow reachO \ s2 \Longrightarrow reachV \ sv1 \Longrightarrow reachV \ sv2 \Longrightarrow$
 $\Delta \ w \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO$
 \Longrightarrow
 $match \ \Delta \ s1 \ s2 \ statA \ sv1 \ sv2 \ statO$

shows *unwindCond* Δ

using *assms unfolding unwindCond-def by auto*

3.3 The soundness of unwinding

definition $\psi \ s1 \ tr1 \ s2 \ tr2 \ statO \ sv1 \ trv1 \ sv2 \ trv2 \equiv$

$trv1 \neq [] \wedge trv2 \neq [] \wedge$
 $Van.validFromS \ sv1 \ trv1 \wedge$

$Van.validFromS\ sv2\ trv2 \wedge$
 $(finalV\ (lastt\ sv1\ trv1) \longleftrightarrow finalO\ (lastt\ s1\ tr1)) \wedge (finalV\ (lastt\ sv2\ trv2) \longleftrightarrow$
 $finalO\ (lastt\ s2\ tr2)) \wedge$
 $Van.S\ trv1 = Opt.S\ tr1 \wedge Van.S\ trv2 = Opt.S\ tr2 \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $(statO = Eq \wedge Opt.O\ tr1 \neq Opt.O\ tr2 \longrightarrow Van.O\ trv1 \neq Van.O\ trv2)$

lemma ψ -completedFrom: $completedFromO\ s1\ tr1 \implies completedFromO\ s2\ tr2 \implies$

$\psi\ s1\ tr1\ s2\ tr2\ statO\ sv1\ trv1\ sv2\ trv2$
 $\implies completedFromV\ sv1\ trv1 \wedge completedFromV\ sv2\ trv2$

unfolding ψ -def $Opt.completedFrom$ -def $Van.completedFrom$ -def $lastt$ -def
by *presburger*

lemma $completedFromO$ -lastt: $completedFromO\ s1\ tr1 \implies finalO\ (lastt\ s1\ tr1)$

unfolding $Opt.completedFrom$ -def $lastt$ -def **by** *auto*

lemma *rsecure-strong*:

assumes

$\bigwedge s1\ tr1\ s2\ tr2.$

$istateO\ s1 \wedge Opt.validFromS\ s1\ tr1 \wedge completedFromO\ s1\ tr1 \wedge$

$istateO\ s2 \wedge Opt.validFromS\ s2\ tr2 \wedge completedFromO\ s2\ tr2 \wedge$

$Opt.A\ tr1 = Opt.A\ tr2 \wedge (isIntO\ s1 \wedge isIntO\ s2 \longrightarrow getActO\ s1 = getActO\ s2)$

\implies

$\exists sv1\ trv1\ sv2\ trv2.$

$istateV\ sv1 \wedge istateV\ sv2 \wedge corrState\ sv1\ s1 \wedge corrState\ sv2\ s2 \wedge$

$\psi\ s1\ tr1\ s2\ tr2\ Eq\ sv1\ trv1\ sv2\ trv2$

shows *rsecure*

unfolding *rsecure-def3* **apply** *clarify*

subgoal for $s1\ tr1\ s2\ tr2$

using *assms*[of $s1\ tr1\ s2\ tr2$]

using ψ -completedFrom ψ -def $completedFromO$ -lastt **by** (*metis* (*full-types*))

.

proposition *unwindCond-ex- ψ* :

assumes *unwind*: $unwindCond\ \Delta$

and Δ : $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$ **and** *stat*: $(statA = Diff \longrightarrow statO = Diff)$

and *v*: $Opt.validFromS\ s1\ tr1\ Opt.completedFrom\ s1\ tr1\ Opt.validFromS\ s2\ tr2$
 $Opt.completedFrom\ s2\ tr2$

and *tr14*: $Opt.A\ tr1 = Opt.A\ tr2$

and *r*: $reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$

shows $\exists trv1\ trv2. \psi\ s1\ tr1\ s2\ tr2\ statO\ sv1\ trv1\ sv2\ trv2$

using *assms*(2-)

proof(*induction* $length\ tr1 + length\ tr2\ w$)

```

    arbitrary: s1 s2 statA sv1 sv2 statO tr1 tr2 rule: less2-induct')
  case (less w tr1 tr2 s1 s2 statA sv1 sv2 statO)
  note ok = ⟨statA = Diff → statO = Diff⟩
  note Δ = ⟨Δ w s1 s2 statA sv1 sv2 statO⟩
  note A34 = ⟨Opt.A tr1 = Opt.A tr2⟩
  note r34 = less.prem(8,9) note r12 = less.prem(10,11)
  note r = r34 r12
  note r3 = r34(1) note r4 = r34(2) note r1 = r12(1) note r2 = r12(2)

  have i34: statA = Eq → isIntO s1 = isIntO s2
  and f34: finalO s1 = finalO s2 ∧ finalV sv1 = finalO s1 ∧ finalV sv2 = finalO
s2
  using Δ unwind[unfolded unwindCond-def] r by auto

  have proact-match: (∃ v < w. proact Δ v s1 s2 statA sv1 sv2 statO) ∨ match Δ s1
s2 statA sv1 sv2 statO
  using Δ unwind[unfolded unwindCond-def] r by auto
  show ?case using proact-match proof safe
  fix v assume v: v < w
  assume proact Δ v s1 s2 statA sv1 sv2 statO
  thus ?thesis unfolding proact-def proof safe
  assume sv1: ¬ isSecV sv1 ¬ isIntV sv1 and move-1 Δ v s1 s2 statA sv1 sv2
statO
  then obtain sv1'
  where 0: validTransV (sv1,sv1')
  and Δ: Δ v s1 s2 statA sv1' sv2 statO
  unfolding move-1-def by auto
  have r1': reachV sv1' using r1 0 by (metis Van.reach.Step fst-conv snd-conv)
  obtain trv1 trv2 where ψ: ψ s1 tr1 s2 tr2 statO sv1' trv1 sv2 trv2
  using less(2)[OF v, of tr1 tr2 s1 s2 statA sv1' sv2 statO, simplified, OF Δ
ok - - - - r34 r1' r2]
  using A34 less.prem(3-6) by blast
  show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
  using ψ ok 0 sv1 unfolding ψ-def Van.completedFrom-def by auto
  next
  assume sv2: ¬ isSecV sv2 ¬ isIntV sv2 and move-2 Δ v s1 s2 statA sv1 sv2
statO
  then obtain sv2'
  where 0: validTransV (sv2,sv2')
  and Δ: Δ v s1 s2 statA sv1 sv2' statO
  unfolding move-2-def by auto
  have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
  obtain trv1 trv2 where ψ: ψ s1 tr1 s2 tr2 statO sv1 trv1 sv2' trv2
  using less(2)[OF v, of tr1 tr2 s1 s2 statA sv1 sv2' statO, simplified, OF Δ
ok - - - - r34 r1 r2']
  using A34 less.prem(3-6) by blast
  show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
  using ψ ok 0 sv2 unfolding ψ-def Van.completedFrom-def by auto
  next

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assume sv12:  $\neg$  isSecV sv1  $\neg$  isSecV sv2 Van.eqAct sv1 sv2
and move-12  $\Delta$  v s1 s2 statA sv1 sv2 statO
then obtain sv1' sv2' statO'
where 0: statO' = sstatO' statO sv1 sv2
validTransV (sv1,sv1')  $\neg$  isSecV sv1
validTransV (sv2,sv2')  $\neg$  isSecV sv2
Van.eqAct sv1 sv2
and  $\Delta$ :  $\Delta$  v s1 s2 statA sv1' sv2' statO'
unfolding move-12-def by auto
have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
fst-conv snd-conv)+
have ok': statA = Diff  $\longrightarrow$  statO' = Diff using ok 0 unfolding sstatO'-def
by (cases statO, auto)
obtain trv1 trv2 where  $\psi$ :  $\psi$  s1 tr1 s2 tr2 statO' sv1' trv1 sv2' trv2
using less(2)[OF v, of tr1 tr2 s1 s2 statA sv1' sv2' statO', simplified, OF  $\Delta$ 
ok' - - - - r34 r12']
using A34 less.prem(3-6) by blast
show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
using  $\psi$  ok' 0 sv12 unfolding  $\psi$ -def sstatO'-def Van.completedFrom-def
using Van.A.Cons-unfold Van.eqAct-def completedFromO-lastt less.prem(4)
less.prem(6) by auto
qed
next
assume m: match  $\Delta$  s1 s2 statA sv1 sv2 statO
show ?thesis
proof(cases length tr1  $\leq$  Suc 0)
case True note tr1 = True
hence tr1 = []  $\vee$  tr1 = [s1]
by (metis Opt.validFromS-Cons-iff le-0-eq le-SucE length-0-conv length-Suc-conv
less.prem(3))
hence finalO s1 using less(3-6)
using Opt.completed-Cons Opt.completed-Nil by blast
hence f4: finalO s2 using f34 by blast
hence tr2: tr2 = []  $\vee$  tr2 = [s2]
by (metis Opt.final-def Simple-Transition-System.validFromS-Cons-iff less.prem(5)
neq-Nil-conv)
show ?thesis apply(rule exI[of - [sv1]], rule exI[of - [sv2]]) using tr1 tr2
using f4 f34
using completedFromO-lastt less.prem(4)
by (auto simp add: lastt-def  $\psi$ -def)
next
case False
then obtain s13 tr1' where tr1: tr1 = s13 # tr1' and tr1'NE: tr1'  $\neq$  []
by (cases tr1, auto)
have s13[simp]: s13 = s1 using  $\langle$ Opt.validFromS s1 tr1 $\rangle$ 
by (simp add: Opt.validFromS-Cons-iff tr1)
obtain s1' where
trn3: validTransO (s1,s1') and

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$tr1'$: *Opt.validFromS* $s1'$ $tr1'$ **using** \langle *Opt.validFromS* $s1$ $tr1$ \rangle
unfolding $tr1$ $s13$ **by** (*metis* $tr1'$ *NE Simple-Transition-System.validFromS-Cons-iff*)
have $r3'$: *reachO* $s1'$ **using** $r3$ $trn3$ **by** (*metis* *Opt.reach.Step fst-conv*
snd-conv)
have $f3$: \neg *finalO* $s1$ **using** *Opt.final-def* $trn3$ **by** *blast*
hence $f4$: \neg *finalO* $s2$ **using** $f34'$ **by** *blast*
hence $tr2$: \neg *length* $tr2 \leq$ *Suc* 0
by (*metis* (*no-types, opaque-lifting*) *Opt.completed-Cons* *Opt.completed-Nil*
Simple-Transition-System.validFromS-Cons-iff *Suc-n-not-le-n* *bot-nat-0.extremum*
le-Suc-eq *length-Cons* *less.premis(5)* *less.premis(6)* *list.exhaust* *order-antisym-conv*)
then obtain $s24$ $tr2'$ **where** $tr2$: $tr2 = s24 \#$ $tr2'$ **and** $tr2'$ *NE*: $tr2' \neq []$
by (*cases* $tr2$, *auto*)
have $s24$ [*simp*]: $s24 = s2$ **using** \langle *Opt.validFromS* $s2$ $tr2$ \rangle
by (*simp* *add*: *Opt.validFromS-Cons-iff* $tr2$)
obtain $s2'$ **where**
 $trn4$: *validTransO* ($s2, s2'$) \vee ($s2 = s2' \wedge$ $tr2' = []$) **and**
 $tr2'$: *Opt.validFromS* $s2'$ $tr2'$ **using** \langle *Opt.validFromS* $s2$ $tr2$ \rangle
unfolding $tr2$ $s24$ **using** *Opt.validFromS-Cons-iff* **by** *auto*
have $r34'$: *reachO* $s1'$ *reachO* $s2'$
using $r3$ $trn3$ $r4$ $trn4$ **by** (*metis* *Opt.reach.Step fst-conv snd-conv*)
note $r3' = r34'(1)$ **note** $r4' = r34'(2)$
define $statA'$ **where** $statA'$: $statA' = sstatA'$ $statA$ $s1$ $s2$
have \neg *isIntO* $s1 \vee$ \neg *isIntO* $s2 \vee$ (*isIntO* $s1 \wedge$ *isIntO* $s2$)
by *auto*
thus *?thesis*
proof *safe*
assume $isAO3$: \neg *isIntO* $s1$
have $O33'$: *Opt.O* $tr1 =$ *Opt.O* $tr1'$ *Opt.A* $tr1 =$ *Opt.A* $tr1'$
using $isAO3$ **unfolding** $tr1$ **by** *auto*
have $A34'$: *Opt.A* $tr1' =$ *Opt.A* $tr2$
using $A34$ $O33'(2)$ **by** *auto*
have m : *match1* Δ $s1$ $s2$ $statA$ $sv1$ $sv2$ $statO$ **using** m **unfolding** *match-def*
by *auto*
have (\neg *isSecO* $s1 \wedge$ $\Delta \infty$ $s1' s2$ $statA$ $sv1$ $sv2$ $statO$) \vee
(*eqSec* $sv1$ $s1 \wedge$ \neg *isIntV* $sv1 \wedge$ *match1-1* Δ $s1$ $s1' s2$ $statA$ $sv1$ $sv2$
 $statO$) \vee
(*eqSec* $sv1$ $s1 \wedge$ \neg *isSecV* $sv2 \wedge$ *Van.eqAct* $sv1$ $sv2 \wedge$ *match1-12* Δ $s1$
 $s1' s2$ $statA$ $sv1$ $sv2$ $statO$)
using m $isAO3$ $trn3$ *ok* **unfolding** *match1-def* **by** *auto*
thus *?thesis*
proof *safe*
assume \neg *isSecO* $s1$ **and** Δ : $\Delta \infty$ $s1' s2$ $statA$ $sv1$ $sv2$ $statO$
hence $S3$: *Opt.S* $tr1' =$ *Opt.S* $tr1$ **unfolding** $tr1$ **by** *auto*
obtain $trv1$ $trv2$ **where** ψ : ψ $s1$ $tr1' s2$ $tr2$ $statO$ $sv1$ $trv1$ $sv2$ $trv2$
using *less(1)* [*of* $tr1' tr2$, *OF* - Δ - - - - - $r3' r4$ $r12$, *unfolded* $O33'$,
simplified]
using *less.premis* $tr1'$ *ok* $A34'$ $f3$ $f4$ **unfolding** $tr1$ *Opt.completedFrom-def*
by (*auto* *split*: *if-splits* *simp*: ψ -*def* *lastt-def*)
show *?thesis* **apply**(*rule* *exI* [*of* - $trv1$]) **apply**(*rule* *exI* [*of* - $trv2$])

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using  $\psi$  O33' S3 unfolding  $\psi$ -def
using completedFromO-lastt less.premis(4)
by (auto simp add: tr1 tr1'NE)
next
assume trn13: eqSec sv1 s1 and
Atrn1:  $\neg$  isIntV sv1 and match1-1  $\Delta$  s1 s1' s2 statA sv1 sv2 statO
then obtain sv1' where
trn1: validTransV (sv1,sv1') and
 $\Delta$ :  $\Delta \infty$  s1' s2 statA sv1' sv2 statO
unfolding match1-1-def by auto
have r1': reachV sv1' using r1 trn1 by (metis Van.reach.Step fst-conv
snd-conv)
obtain trv1 trv2 where  $\psi$ :  $\psi$  s1 tr1' s2 tr2 statO sv1' trv1 sv2 trv2
using less(1)[of tr1' tr2, OF -  $\Delta$  - - - - - r3' r4 r1' r2, unfolded O33',
simplified]
using less.premis tr1' ok A34' f3 f4 unfolding tr1 tr2 Opt.completedFrom-def

by (auto simp:  $\psi$ -def lastt-def split: if-splits)
show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
using  $\psi$  O33' unfolding tr1 tr2 Van.completedFrom-def
using Van.validFromS-Cons trn1 tr1'NE tr2'NE
using isAO3 ok Atrn1 eqSec-S-Cons trn13
unfolding  $\psi$ -def using completedFromO-lastt less.premis(4) tr1 by auto

next
assume sv2:  $\neg$  isSecV sv2 and trn13: eqSec sv1 s1 and
Atrn12: Van.eqAct sv1 sv2 and match1-12  $\Delta$  s1 s1' s2 statA sv1 sv2 statO
then obtain sv1' sv2' statO' where
statO': statO' = sstatO' statO sv1 sv2 and
trn1: validTransV (sv1,sv1') and
trn2: validTransV (sv2,sv2') and
 $\Delta$ :  $\Delta \infty$  s1' s2 statA sv1' sv2' statO'
unfolding match1-12-def by auto
have r12': reachV sv1' reachV sv2'
using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
obtain trv1 trv2 where  $\psi$ :  $\psi$  s1' tr1' s2 tr2 statO' sv1' trv1 sv2' trv2
using less(1)[of tr1' tr2, OF -  $\Delta$  - - - - - r3' r4 r12', unfolded O33',
simplified]
using less.premis tr1' ok A34' f3 f4 unfolding tr1 tr2 Opt.completedFrom-def
statO' sstatO'-def
by auto presburger+
show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
using  $\psi$  O33' tr1'NE tr2'NE sv2
using Van.validFromS-Cons trn1 trn2
using isAO3 ok Atrn12 eqSec-S-Cons trn13 f3 f34 s13
unfolding  $\psi$ -def tr1 Van.completedFrom-def Van.eqAct-def statO' sstatO'-def
using Van.A.Cons-unfold tr1' trn3 by auto
qed

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next
  assume isAO4:  $\neg$  isIntO s2
  have O44':  $Opt.O\ tr2 = Opt.O\ tr2'\ Opt.A\ tr2 = Opt.A\ tr2'$ 
  using isAO4 unfolding tr2 by auto
  have A34':  $Opt.A\ tr1 = Opt.A\ tr2'$ 
  using A34 O44'(2) by auto
  have m: match2  $\Delta\ s1\ s2\ statA\ sv1\ sv2\ statO$  using m unfolding match-def
by auto
  have ( $\neg$  isSecO s2  $\wedge\ \Delta\ \infty\ s1\ s2'\ statA\ sv1\ sv2\ statO$ )  $\vee$ 
    (eqSec sv2 s2  $\wedge\ \neg$  isIntV sv2  $\wedge\ match2-1\ \Delta\ s1\ s2\ s2'\ statA\ sv1\ sv2$ 
statO)  $\vee$ 
    ( $\neg$  isSecV sv1  $\wedge\ eqSec\ sv2\ s2\ \wedge\ Van.eqAct\ sv1\ sv2\ \wedge\ match2-12\ \Delta\ s1$ 
s2 s2' statA sv1 sv2 statO)
  using m isAO4 trn4 ok tr2'NE unfolding match2-def by auto
  thus ?thesis
  proof safe
    assume  $\neg$  isSecO s2 and  $\Delta: \Delta\ \infty\ s1\ s2'\ statA\ sv1\ sv2\ statO$ 
    hence S4:  $Opt.S\ tr2' = Opt.S\ tr2$  unfolding tr2 by auto
    obtain trv1 trv2 where  $\psi: \psi\ s1\ tr1\ s2'\ tr2'\ statO\ sv1\ trv1\ sv2\ trv2$ 
    using less(1)[of tr1 tr2', OF -  $\Delta$  - - - - - r3 r4', simplified]
    using less.premis tr2' ok A34' tr1'NE tr2'NE unfolding tr1 tr2 Opt.completedFrom-def
by (cases isIntO s2, auto)
    show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
    using  $\psi\ O44'\ S4$  unfolding  $\psi$ -def
    using completedFromO-lastt less.premis(6)
    unfolding Opt.completedFrom-def using tr2 tr2'NE by auto
  next
    assume trn24: eqSec sv2 s2 and
    Atrn2:  $\neg$  isIntV sv2 and match2-1  $\Delta\ s1\ s2\ s2'\ statA\ sv1\ sv2\ statO$ 
    then obtain sv2' where trn2: validTransV (sv2,sv2') and
     $\Delta: \Delta\ \infty\ s1\ s2'\ statA\ sv1\ sv2'\ statO$ 
    unfolding match2-1-def by auto
    have r2': reachV sv2' using r2 trn2 by (metis Van.reach.Step fst-conv
snd-conv)
    obtain trv1 trv2 where  $\psi: \psi\ s1\ tr1\ s2'\ tr2'\ statO\ sv1\ trv1\ sv2'\ trv2$ 
    using less(1)[of tr1 tr2', OF -  $\Delta$  - - - - - r3 r4' r1 r2', simplified]
    using less.premis tr2' ok A34' tr1'NE tr2'NE unfolding tr1 tr2 Opt.completedFrom-def
by (cases isIntO s2, auto)
    show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
    using  $\psi\ tr1'NE\ tr2'NE$ 
    using Van.validFromS-Cons trn2
    using isAO4 ok Atrn2 eqSec-S-Cons trn24
    unfolding  $\psi$ -def tr1 tr2 s13 s24 Van.completedFrom-def lastt-def by auto
  next
    assume sv1:  $\neg$  isSecV sv1 and trn24: eqSec sv2 s2 and
    Atrn12: Van.eqAct sv1 sv2 and match2-12  $\Delta\ s1\ s2\ s2'\ statA\ sv1\ sv2$ 
statO
    then obtain sv1' sv2' statO' where
    statO':  $statO' = sstatO'\ statO\ sv1\ sv2$  and

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    trn1: validTransV (sv1,sv1') and
    trn2: validTransV (sv2,sv2') and
    Δ: Δ ∞ s1 s2' statA sv1' sv2' statO'
    unfolding match2-12-def by auto
    have r12': reachV sv1' reachV sv2'
    using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
    obtain trv1 trv2 where ψ: ψ s1 tr1 s2' tr2' statO' sv1' trv1 sv2' trv2
    using less(1)[of tr1 tr2', OF - Δ - - - - - r3 r4' r12', simplified]
    using less.premis tr2' ok A34' tr1'NE tr2'NE unfolding tr1 tr2 Opt.completedFrom-def
statO' sstatO'-def
    by (cases isIntO s2, auto)
    show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
    using ψ O44' tr1'NE tr2'NE sv1
    using Van.validFromS-Cons trn1 trn2
    using isAO4 ok Atrn12 eqSec-S-Cons trn24
    unfolding ψ-def tr2 tr1'NE Van.completedFrom-def Van.eqAct-def
statO' sstatO'-def
    using Van.A.Cons-unfold tr2' trn4 by auto
qed
next
assume isAO34: isIntO s1 isIntO s2
have A34': getActO s1 = getActO s2 Opt.A tr1' = Opt.A tr2'
using A34 isAO34 tr1'NE tr2'NE unfolding tr1 tr2 by auto
have O33': Opt.O tr1 = getObsO s1 # Opt.O tr1' and
    O44': Opt.O tr2 = getObsO s2 # Opt.O tr2'
using isAO34 tr1'NE tr2'NE unfolding tr1 s13 tr2 s24 by auto
have m: match12 Δ s1 s2 statA sv1 sv2 statO using m unfolding statA'
match-def by auto
have trn34: getObsO s1 = getObsO s2 ∨ statA' = Diff
using isAO34 unfolding statA' sstatA'-def by (cases statA,auto)
have (¬ isSecO s1 ∧ ¬ isSecO s2 ∧ (statA = statA' ∨ statO = Diff) ∧ Δ
∞ s1' s2' statA' sv1 sv2 statO)
  ∨
  (¬ isSecO s2 ∧ eqSec sv1 s1 ∧ ¬ isIntV sv1 ∧
  (statA = statA' ∨ statO = Diff) ∧
  match12-1 Δ s1' s2' statA' sv1 sv2 statO)
  ∨
  (¬ isSecO s1 ∧ eqSec sv2 s2 ∧ ¬ isIntV sv2 ∧
  (statA = statA' ∨ statO = Diff) ∧
  match12-2 Δ s1' s2' statA' sv1 sv2 statO)
  ∨
  (eqSec sv1 s1 ∧ eqSec sv2 s2 ∧ Van.eqAct sv1 sv2 ∧
  match12-12 Δ s1' s2' statA' sv1 sv2 statO)
(is ?K1 ∨ ?K2 ∨ ?K3 ∨ ?K4)
using m[unfolded match12-def, rule-format, of s1' s2']
isAO34 A34' trn3 trn4 tr1'NE tr2'NE
unfolding s13 s24 trn34 statA' Opt.eqAct-def sstatA'-def by auto
thus ?thesis proof (elim disjE)

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assume  $K1$ : ? $K1$  hence  $\Delta$ :  $\Delta \infty s1' s2' statA' sv1 sv2 statO$  by simp
have  $ok'$ : ( $statA' = Diff \longrightarrow statO = Diff$ )
using  $ok$   $K1$  unfolding  $statA'$  using  $isAO34'$  by auto
obtain  $trv1$   $trv2$  where  $\psi$ :  $\psi s1' tr1' s2' tr2' statO sv1 trv1 sv2 trv2$ 
using  $less(1)[of\ tr1'\ tr2',\ OF - \Delta - - - - - r34'\ r12,\ simplified]$ 
using  $less.premis\ tr1'\ tr2'\ ok'\ A34'\ isAO34'\ tr1'NE\ tr2'NE$  unfolding  $tr1$ 
 $tr2$  Opt.comPLETEDFrom-def by auto
show ?thesis apply(rule  $exI[of - trv1]$ ) apply(rule  $exI[of - trv2]$ )
using  $\psi$   $trn34'$   $O33'$   $O44'$   $K1$   $ok$  unfolding  $\psi$ -def  $tr1$   $tr2$ 
using completedFromO-lastt  $less.premis(4,6)$ 
unfolding Opt.comPLETEDFrom-def using  $tr1$   $tr2$   $tr1'NE$   $tr2'NE$  by auto
next
assume  $K2$ : ? $K2$ 
then obtain  $sv1'$  where
 $trn1$ : validTransV ( $sv1,sv1'$ ) and
 $trn13$ : eqSec  $sv1$   $s1$  and
 $Atrn1$ :  $\neg isIntV\ sv1$  and  $ok'$ : ( $statA' = statA \vee statO = Diff$ ) and
 $\Delta$ :  $\Delta \infty s1' s2' statA' sv1' sv2 statO$ 
unfolding match12-1-def by auto
have  $r1'$ : reachV  $sv1'$  using  $r1$   $trn1$  by (metis Van.reach.Step fst-conv
snd-conv)
obtain  $trv1$   $trv2$  where  $\psi$ :  $\psi s1' tr1' s2' tr2' statO sv1' trv1 sv2 trv2$ 
using  $less(1)[of\ tr1'\ tr2',\ OF - \Delta - - - - - r34'\ r1'\ r2,\ simplified]$ 
using  $less.premis\ tr1'\ tr2'\ ok'\ A34'\ tr1'NE\ tr2'NE$  unfolding  $tr1$   $tr2$ 
Opt.comPLETEDFrom-def by auto
show ?thesis apply(rule  $exI[of - sv1 \# trv1]$ ) apply(rule  $exI[of - trv2]$ )
using  $\psi$   $O33'$   $O44'$   $tr1'NE$   $tr2'NE$  unfolding  $tr1$   $tr2$ 
using Van.validFromS-Cons  $trn1$   $ok$ 
using  $K2$   $ok'$   $Atrn1$  eqSec-S-Cons  $trn13$   $trn34'$ 
unfolding  $statA'$  Van.comPLETEDFrom-def eqSec-def
using  $s13$   $tr1$   $tr1'$   $tr2'$   $trn3$   $trn4$ 
by simp (smt (verit, ccfv-SIG) Opt.S.simps(2) Simple-Transition-System.lastt-Cons
Van.A.Cons-unfold Van.O.Cons-unfold  $\psi$ -def list.simps(3) status.simps(1))
next
assume  $K3$ : ? $K3$ 
then obtain  $sv2'$  where
 $trn2$ : validTransV ( $sv2,sv2'$ ) and
 $trn24$ : eqSec  $sv2$   $s2$  and
 $Atrn2$ :  $\neg isIntV\ sv2$  and  $ok'$ : ( $statA' = statA \vee statO = Diff$ ) and
 $\Delta$ :  $\Delta \infty s1' s2' statA' sv1 sv2' statO$ 
unfolding match12-2-def by auto
have  $r2'$ : reachV  $sv2'$  using  $r2$   $trn2$  by (metis Van.reach.Step fst-conv
snd-conv)
obtain  $trv1$   $trv2$  where  $\psi$ :  $\psi s1' tr1' s2' tr2' statO sv1 trv1 sv2' trv2$ 
using  $less(1)[of\ tr1'\ tr2',\ OF - \Delta - - - - - r34'\ r1\ r2',\ simplified]$ 
using  $less.premis\ tr1'\ tr2'\ ok'\ A34'\ tr1'NE\ tr2'NE$  unfolding  $tr1$   $tr2$ 
Opt.comPLETEDFrom-def by auto
show ?thesis apply(rule  $exI[of - trv1]$ ) apply(rule  $exI[of - sv2 \# trv2]$ )

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using  $\psi$  O33' O44' tr1'NE tr2'NE unfolding  $\psi$ -def tr1 tr2
using Van.validFromS-Cons trn2 ok
using K3 ok' Atrn2 eqSec-S-Cons trn24 trn34
unfolding statA' Van.completedFrom-def
by simp (metis last.simps lastt-def list.simps(3) status.distinct(2))
next
assume K4: ?K4
then obtain sv1' sv2' statO' where 0:
  statO' = sstatO' statO sv1 sv2
  validTransV (sv1,sv1')
  eqSec sv1 s1
  validTransV (sv2,sv2')
  eqSec sv2 s2
  Van.eqAct sv1 sv2
  and ok': statA' = Diff  $\longrightarrow$  statO' = Diff and  $\Delta$ :  $\Delta \infty s1' s2' statA'$ 
sv1' sv2' statO'
  unfolding match12-12-def by auto
  have r12': reachV sv1' reachV sv2' using r1 r2 0
  by (metis Van.reach.Step fst-conv snd-conv)+
  obtain trv1 trv2 where  $\psi$ :  $\psi s1' tr1' s2' tr2' statO' sv1' trv1 sv2' trv2$ 
  using less(1)[of tr1' tr2', OF -  $\Delta$  - - - - - r34' r12', simplified]
  using less.premis tr1' tr2' ok' A34' tr1'NE tr2'NE unfolding tr1 tr2
Opt.completedFrom-def by auto
  show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
  using trn34
  using  $\psi$  O33' O44' isAO34 tr1'NE tr2'NE unfolding  $\psi$ -def tr1 tr2
  using Van.validFromS-Cons 0
  using K4 eqSec-S-Cons
  unfolding statA' Van.eqAct-def Van.completedFrom-def match12-12-def
sstatO'-def
  by (smt (z3) Simple-Transition-System.lastt-Cons Van.A.Cons-unfold
Van.O.Cons-unfold
  completedFromO-lastt f3 f34 lastt-Nil less.premis(4) less.premis(6) list.inject
s13
  s24 status.simps(1) tr1 tr1' tr2 tr2' trn3 trn4 updStat.simps(1) upd-
Stat.simps(3))
qed
qed
qed
qed
qed

```

theorem *unwind-rsecure*:

```

assumes init: initCond  $\Delta$ 
  and unwind: unwindCond  $\Delta$ 
shows rsecure
apply (rule rsecure-strong)
apply (elim conjE)

```

```

subgoal for  $s1\ tr1\ s2\ tr2$ 
  using init unfolding initCond-def
  apply (erule-tac allE[of - s1])
  apply (elim allE[of - s2] conjE)
  apply (elim impE[of ⟨istateO s1 ∧ istateO s2⟩] exE conjE)
  subgoal by clarify
  subgoal for  $sv1\ sv2$ 
    using unwind apply (drule-tac unwindCond-ex-ψ, blast+)
    subgoal by (rule Transition-System.reach.Istate)
    subgoal by (rule Transition-System.reach.Istate)
    subgoal by (rule Transition-System.reach.Istate)
    subgoal by (rule Transition-System.reach.Istate)
    apply (elim exE)
    subgoal for  $trv1\ trv2$ 
      apply (rule exI[of - sv1], rule exI[of - trv1], rule exI[of - sv2], rule exI[of -
by clarify
    .
  .
.

```

3.4 Compositional unwinding

We allow networks of unwinding relations that unwind into each other, which offer a compositional alternative to monolithic unwinding.

definition *unwindIntoCond* ::

$$\begin{aligned}
& (enat \Rightarrow 'stateO \Rightarrow 'stateO \Rightarrow status \Rightarrow 'stateV \Rightarrow 'stateV \Rightarrow status \Rightarrow bool) \Rightarrow \\
& (enat \Rightarrow 'stateO \Rightarrow 'stateO \Rightarrow status \Rightarrow 'stateV \Rightarrow 'stateV \Rightarrow status \Rightarrow bool) \\
& \Rightarrow bool
\end{aligned}$$

where

$$\begin{aligned}
& unwindIntoCond\ \Delta\ \Delta' \equiv \forall w\ s1\ s2\ statA\ sv1\ sv2\ statO. \\
& reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge \\
& \Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \longrightarrow \\
& (finalO\ s1 \longleftrightarrow finalO\ s2) \wedge (finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO \\
& s2) \\
& \wedge \\
& (statA = Eq \longrightarrow (isIntO\ s1 \longleftrightarrow isIntO\ s2)) \\
& \wedge \\
& ((\exists v < w. proact\ \Delta'\ v\ s1\ s2\ statA\ sv1\ sv2\ statO) \\
& \vee \\
& match\ \Delta'\ s1\ s2\ statA\ sv1\ sv2\ statO)
\end{aligned}$$

lemma *unwindIntoCond-simpleI*:

assumes

$$\begin{aligned}
& \wedge w\ s1\ s2\ statA\ sv1\ sv2\ statO. \\
& reachO\ s1 \Longrightarrow reachO\ s2 \Longrightarrow reachV\ sv1 \Longrightarrow reachV\ sv2 \Longrightarrow \\
& \Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \\
& \Longrightarrow
\end{aligned}$$

$(finalO\ s1 \longleftrightarrow finalO\ s2) \wedge (finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO\ s2)$
and
 $\bigwedge w\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$
 $statA = Eq$
 \implies
 $isIntO\ s1 \longleftrightarrow isIntO\ s2$
 $\bigwedge w\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$
 \implies
 $match\ \Delta'\ s1\ s2\ statA\ sv1\ sv2\ statO$
shows *unwindIntoCond* $\Delta\ \Delta'$
using *assms unfolding unwindIntoCond-def by auto*

lemma *unwindIntoCond-simpleI2*:

assumes

$\bigwedge w\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$
 \implies
 $(finalO\ s1 \longleftrightarrow finalO\ s2) \wedge (finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO\ s2)$

and

$\bigwedge w\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$
 $statA = Eq$
 \implies
 $isIntO\ s1 \longleftrightarrow isIntO\ s2$

and

$\bigwedge w\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$
 \implies
 $(\exists v < w. proact\ \Delta'\ v\ s1\ s2\ statA\ sv1\ sv2\ statO)$

shows *unwindIntoCond* $\Delta\ \Delta'$

using *assms unfolding unwindIntoCond-def by auto*

lemma *unwindIntoCond-simpleIB*:

assumes

$\bigwedge w\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$
 \implies
 $(finalO\ s1 \longleftrightarrow finalO\ s2) \wedge (finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO\ s2)$

and
 $\bigwedge w s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w s1 s2 statA sv1 sv2 statO \implies$
 $statA = Eq$
 \implies
 $isIntO s1 \longleftrightarrow isIntO s2$
and
 $\bigwedge w s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w s1 s2 statA sv1 sv2 statO$
 \implies
 $(\exists v < w. proact \Delta' v s1 s2 statA sv1 sv2 statO) \vee match \Delta' s1 s2 statA sv1 sv2$
 $statO$
shows *unwindIntoCond* $\Delta \Delta'$
using *assms unfolding unwindIntoCond-def by auto*

theorem *distrib-unwind-rsecure*:
assumes $m: 0 < m$ **and** $nxt: \bigwedge i. i < (m::nat) \implies nxt i \subseteq \{0..<m\}$
and *init*: *initCond* $(\Delta s 0)$
and *step*: $\bigwedge i. i < m \implies$
 $unwindIntoCond (\Delta s i) (\lambda w s1 s2 statA sv1 sv2 statO.$
 $\exists j \in nxt i. \Delta s j w s1 s2 statA sv1 sv2 statO)$
shows *rsecure*
proof –
define Δ **where** $D: \Delta \equiv \lambda w s1 s2 statA sv1 sv2 statO. \exists i < m. \Delta s i w s1 s2$
 $statA sv1 sv2 statO$
have $i: initCond \Delta$
using *init m unfolding initCond-def D by meson*
have $c: unwindCond \Delta$ **unfolding** *unwindCond-def* **apply**(*intro allI impI allI*)
apply(*subst (asm) D*) **apply** (*elim exE conjE*)
subgoal for $w s1 s2 statA sv1 sv2 statO i$
apply(*frule step*) **unfolding** *unwindIntoCond-def*
apply(*erule allE[of - w]*)
apply(*erule allE[of - s1]*) **apply**(*erule allE[of - s2]*) **apply**(*erule allE[of -*
 $statA]$)
apply(*erule allE[of - sv1]*) **apply**(*erule allE[of - sv2]*) **apply**(*erule allE[of -*
 $statO]$)
apply *simp* **apply**(*elim conjE*)
apply(*erule disjE*)
subgoal **apply**(*rule disjI1*)
subgoal **apply**(*elim exE conjE*) **subgoal for** v
apply(*rule exI[of - v], simp*)
apply(*rule proact-mono[of $\lambda w s1 s2 statA sv1 sv2 statO. \exists j \in nxt i. \Delta s j w$*
 $s1 s2 statA sv1 sv2 statO]$)
subgoal **unfolding** *le-fun-def D by simp (meson atLeastLessThan-iff nxt*
 $subsetD)$
subgoal
subgoal **apply**(*rule disjI2*)

apply(rule match-mono[of $\lambda w s1 s2 statA sv1 sv2 statO. \exists j \in next\ i. \Delta s\ j\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$])
subgoal unfolding le-fun-def *D* by simp (meson atLeastLessThan-iff next subsetD)
subgoal
show ?thesis using unwind-rsecure[OF *i c*].
qed

corollary linear-unwind-rsecure:
assumes *init*: initCond ($\Delta s\ 0$)
and *step*: ($\bigwedge i. i < m \implies$
unwindIntoCond ($\Delta s\ i$) ($\lambda w s1 s2 statA sv1 sv2 statO.$
 $\Delta s\ i\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \vee$
 $\Delta s\ (Suc\ i)\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$)
and *finish*: unwindIntoCond ($\Delta s\ m$) ($\Delta s\ m$)
shows rsecure
apply(rule distrib-unwind-rsecure[of *Suc m* $\lambda i. \text{if } i < m \text{ then } \{i, Suc\ i\} \text{ else } \{i\}$ $\Delta s,$
OF - - init])
using *step finish*
by (auto simp: less-Suc-eq-le)

definition oor where
 $oor\ \Delta\ \Delta_2 \equiv \lambda w s1 s2 statA sv1 sv2 statO.$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \vee \Delta_2\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$

lemma oorI1:
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \implies oor\ \Delta\ \Delta_2\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$
unfolding oor-def by simp

lemma oorI2:
 $\Delta_2\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \implies oor\ \Delta\ \Delta_2\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$
unfolding oor-def by simp

definition oor3 where
 $oor3\ \Delta\ \Delta_2\ \Delta_3 \equiv \lambda w s1 s2 statA sv1 sv2 statO.$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \vee \Delta_2\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \vee$
 $\Delta_3\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$

lemma oor3I1:
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \implies oor3\ \Delta\ \Delta_2\ \Delta_3\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$
unfolding oor3-def by simp

lemma oor3I2:
 $\Delta_2\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \implies oor3\ \Delta\ \Delta_2\ \Delta_3\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$
unfolding oor3-def by simp

lemma oor3I3:

$\Delta_3 w s1 s2 statA sv1 sv2 statO \implies oor3 \Delta \Delta_2 \Delta_3 w s1 s2 statA sv1 sv2 statO$
unfolding oor3-def by simp

definition oor4 where

$oor4 \Delta \Delta_2 \Delta_3 \Delta_4 \equiv \lambda w s1 s2 statA sv1 sv2 statO.$
 $\Delta w s1 s2 statA sv1 sv2 statO \vee \Delta_2 w s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_3 w s1 s2 statA sv1 sv2 statO \vee \Delta_4 w s1 s2 statA sv1 sv2 statO$

lemma oor4I1:

$\Delta w s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w s1 s2 statA sv1 sv2 statO$
unfolding oor4-def by simp

lemma oor4I2:

$\Delta_2 w s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w s1 s2 statA sv1 sv2 statO$
unfolding oor4-def by simp

lemma oor4I3:

$\Delta_3 w s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w s1 s2 statA sv1 sv2 statO$
unfolding oor4-def by simp

lemma oor4I4:

$\Delta_4 w s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w s1 s2 statA sv1 sv2 statO$
unfolding oor4-def by simp

definition oor5 where

$oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \equiv \lambda w s1 s2 statA sv1 sv2 statO.$
 $\Delta w s1 s2 statA sv1 sv2 statO \vee \Delta_2 w s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_3 w s1 s2 statA sv1 sv2 statO \vee \Delta_4 w s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_5 w s1 s2 statA sv1 sv2 statO$

lemma oor5I1:

$\Delta w s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w s1 s2 statA sv1 sv2 statO$
unfolding oor5-def by simp

lemma oor5I2:

$\Delta_2 w s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w s1 s2 statA sv1 sv2 statO$
unfolding oor5-def by simp

lemma oor5I3:

$\Delta_3 w s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w s1 s2 statA sv1 sv2 statO$
unfolding oor5-def by simp

lemma oor5I4:

$\Delta_4 w s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w s1 s2 statA sv1 sv2 statO$
unfolding oor5-def by simp

lemma *oor5I5*:

$\Delta_5 w s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def* by *simp*

definition *oor6* **where**

$oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 \equiv \lambda w s1 s2 statA sv1 sv2 statO.$

$\Delta w s1 s2 statA sv1 sv2 statO \vee \Delta_2 w s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_3 w s1 s2 statA sv1 sv2 statO \vee \Delta_4 w s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_5 w s1 s2 statA sv1 sv2 statO \vee \Delta_6 w s1 s2 statA sv1 sv2 statO$

lemma *oor6I1*:

$\Delta w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def* by *simp*

lemma *oor6I2*:

$\Delta_2 w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def* by *simp*

lemma *oor6I3*:

$\Delta_3 w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def* by *simp*

lemma *oor6I4*:

$\Delta_4 w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def* by *simp*

lemma *oor6I5*:

$\Delta_5 w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def* by *simp*

lemma *oor6I6*:

$\Delta_6 w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def* by *simp*

lemma *unwind-rsecure-foo*:

assumes *init*: *initCond* Δ_0

and *step0*: *unwindIntoCond* $\Delta_0 \Delta NN$

and *stepNN*: *unwindIntoCond* ΔNN (*oor5* $\Delta NN \Delta SN \Delta NS \Delta SS \Delta nonspec$)

```

and stepNS: unwindIntoCond  $\Delta NS$  (oor4  $\Delta NN$   $\Delta SN$   $\Delta NS$   $\Delta SS$ )
and stepSN: unwindIntoCond  $\Delta SN$  (oor4  $\Delta NN$   $\Delta SN$   $\Delta NS$   $\Delta SS$ )
and stepSS: unwindIntoCond  $\Delta SS$  (oor4  $\Delta NN$   $\Delta SN$   $\Delta NS$   $\Delta SS$ )
and stepNonspec: unwindIntoCond  $\Delta nonspec$   $\Delta nonspec$ 
shows rsecure
proof –
define m where m: m  $\equiv$  (6::nat)
define  $\Delta s$  where  $\Delta s$ :  $\Delta s \equiv \lambda i::nat.$ 
  if i = 0 then  $\Delta_0$ 
  else if i = 1 then  $\Delta NN$ 
  else if i = 2 then  $\Delta SN$ 
  else if i = 3 then  $\Delta NS$ 
  else if i = 4 then  $\Delta SS$ 
  else  $\Delta nonspec$ 
define next where next: next  $\equiv \lambda i::nat.$ 
  if i = 0 then {1::nat}
  else if i = 1 then {1,2,3,4,5}
  else if i = 2 then {1,2,3,4}
  else if i = 3 then {1,2,3,4}
  else if i = 4 then {1,2,3,4}
  else {5}
show ?thesis apply(rule distrib-unwind-rsecure[of m next  $\Delta s$ ])
  subgoal unfolding m by auto
  subgoal unfolding next m by auto
  subgoal using init unfolding  $\Delta s$  by auto
  subgoal
    unfolding m next  $\Delta s$ 
    using step0 stepNN stepNS stepSN stepSS stepNonspec
    unfolding oor4-def oor5-def by auto .
qed

```

lemma isIntO-match1: isIntO s1 \implies match1 Δ s1 s2 statA sv1 sv2 statO
unfolding match1-def **by** auto

lemma isIntO-match2: isIntO s2 \implies match2 Δ s1 s2 statA sv1 sv2 statO
unfolding match2-def **by** auto

lemma match1-1-oorI1:
 match1-1 Δ s1 s1' s2 statA sv1 sv2 statO \implies
 match1-1 (oor Δ Δ_2) s1 s1' s2 statA sv1 sv2 statO
apply(rule match1-1-mono) **unfolding** le-fun-def oor-def **by** auto

lemma match1-1-oorI2:
 match1-1 Δ_2 s1 s1' s2 statA sv1 sv2 statO \implies
 match1-1 (oor Δ Δ_2) s1 s1' s2 statA sv1 sv2 statO

apply(rule *match1-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match1-oorI1*:

match1 Δ *s1 s2 statA sv1 sv2 statO* \implies
match1 (*oor* Δ Δ_2) *s1 s2 statA sv1 sv2 statO*

apply(rule *match1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match1-oorI2*:

match1 Δ_2 *s1 s2 statA sv1 sv2 statO* \implies
match1 (*oor* Δ Δ_2) *s1 s2 statA sv1 sv2 statO*

apply(rule *match1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match2-1-oorI1*:

match2-1 Δ *s1 s2 s2' statA sv1 sv2 statO* \implies
match2-1 (*oor* Δ Δ_2) *s1 s2 s2' statA sv1 sv2 statO*

apply(rule *match2-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match2-1-oorI2*:

match2-1 Δ_2 *s1 s2 s2' statA sv1 sv2 statO* \implies
match2-1 (*oor* Δ Δ_2) *s1 s2 s2' statA sv1 sv2 statO*

apply(rule *match2-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match2-oorI1*:

match2 Δ *s1 s2 statA sv1 sv2 statO* \implies
match2 (*oor* Δ Δ_2) *s1 s2 statA sv1 sv2 statO*

apply(rule *match2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match2-oorI2*:

match2 Δ_2 *s1 s2 statA sv1 sv2 statO* \implies
match2 (*oor* Δ Δ_2) *s1 s2 statA sv1 sv2 statO*

apply(rule *match2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-oorI1*:

match12 Δ *s1 s2 statA sv1 sv2 statO* \implies
match12 (*oor* Δ Δ_2) *s1 s2 statA sv1 sv2 statO*

apply(rule *match12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-oorI2*:

match12 Δ_2 *s1 s2 statA sv1 sv2 statO* \implies
match12 (*oor* Δ Δ_2) *s1 s2 statA sv1 sv2 statO*

apply(rule *match12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-1-oorI1*:

match12-1 Δ *s1' s2' statA' sv1 sv2 statO* \implies
match12-1 (*oor* Δ Δ_2) *s1' s2' statA' sv1 sv2 statO*

apply(rule *match12-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-1-oorI2*:

match12-1 Δ_2 *s1' s2' statA' sv1 sv2 statO* \implies

match12-1 (*oor* Δ Δ_2) *s1' s2' statA' sv1 sv2 statO*

apply(rule *match12-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-2-oorI1*:

match12-2 Δ *s2 s2' statA' sv1 sv2 statO* \implies

match12-2 (*oor* Δ Δ_2) *s2 s2' statA' sv1 sv2 statO*

apply(rule *match12-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-2-oorI2*:

match12-2 Δ_2 *s2 s2' statA' sv1 sv2 statO* \implies

match12-2 (*oor* Δ Δ_2) *s2 s2' statA' sv1 sv2 statO*

apply(rule *match12-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-12-oorI1*:

match12-12 Δ *s1' s2' statA' sv1 sv2 statO* \implies

match12-12 (*oor* Δ Δ_2) *s1' s2' statA' sv1 sv2 statO*

apply(rule *match12-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-12-oorI2*:

match12-12 Δ_2 *s1' s2' statA' sv1 sv2 statO* \implies

match12-12 (*oor* Δ Δ_2) *s1' s2' statA' sv1 sv2 statO*

apply(rule *match12-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match-oorI1*:

match Δ *s1 s2 statA sv1 sv2 statO* \implies

match (*oor* Δ Δ_2) *s1 s2 statA sv1 sv2 statO*

apply(rule *match-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match-oorI2*:

match Δ_2 *s1 s2 statA sv1 sv2 statO* \implies

match (*oor* Δ Δ_2) *s1 s2 statA sv1 sv2 statO*

apply(rule *match-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *proact-oorI1*:

proact Δ *meas s1 s2 statA sv1 sv2 statO* \implies

proact (*oor* Δ Δ_2) *meas s1 s2 statA sv1 sv2 statO*

apply(rule *proact-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *proact-oorI2*:

proact Δ_2 *meas s1 s2 statA sv1 sv2 statO* \implies

proact (*oor* Δ Δ_2) *meas s1 s2 statA sv1 sv2 statO*

apply(rule proact-mono) **unfolding** le-fun-def oor-def **by** auto

lemma move-1-oorI1:

$move-1 \Delta meas s1 s2 statA sv1 sv2 statO \implies$

$move-1 (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$

apply(rule move-1-mono) **unfolding** le-fun-def oor-def **by** auto

lemma move-1-oorI2:

$move-1 \Delta_2 meas s1 s2 statA sv1 sv2 statO \implies$

$move-1 (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$

apply(rule move-1-mono) **unfolding** le-fun-def oor-def **by** auto

lemma move-2-oorI1:

$move-2 \Delta meas s1 s2 statA sv1 sv2 statO \implies$

$move-2 (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$

apply(rule move-2-mono) **unfolding** le-fun-def oor-def **by** auto

lemma move-2-oorI2:

$move-2 \Delta_2 meas s1 s2 statA sv1 sv2 statO \implies$

$move-2 (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$

apply(rule move-2-mono) **unfolding** le-fun-def oor-def **by** auto

lemma move-12-oorI1:

$move-12 \Delta meas s1 s2 statA sv1 sv2 statO \implies$

$move-12 (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$

apply(rule move-12-mono) **unfolding** le-fun-def oor-def **by** auto

lemma move-12-oorI2:

$move-12 \Delta_2 meas s1 s2 statA sv1 sv2 statO \implies$

$move-12 (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$

apply(rule move-12-mono) **unfolding** le-fun-def oor-def **by** auto

end

context Relative-Security-Determ

begin

lemma match1-1-defD: $match1-1 \Delta s1 s1' s2 statA sv1 sv2 statO \iff$

$\neg finalV sv1 \wedge \Delta \infty s1' s2 statA (nextO sv1) sv2 statO$

unfolding match1-1-def validTrans-iff-next **by** simp

lemma match1-12-defD: $match1-12 \Delta s1 s1' s2 statA sv1 sv2 statO \iff$

$\neg finalV sv1 \wedge \neg finalV sv2 \wedge$

$\Delta \infty s1' s2 statA (nextO sv1) (nextO sv2) (sstatO' statO sv1 sv2)$

unfolding match1-12-def validTrans-iff-next **by** simp

lemmas match1-defsD = match1-def match1-1-defD match1-12-defD

lemma *match2-1-defD*: $match2-1 \Delta s1 s2 s2' statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv2 \wedge \Delta \infty s1 s2' statA sv1 (nextO sv2) statO$
unfolding *match2-1-def validTrans-iff-next* **by** *simp*

lemma *match2-12-defD*: $match2-12 \Delta s1 s2 s2' statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv1 \wedge \neg finalV sv2 \wedge \Delta \infty s1 s2' statA (nextO sv1) (nextO sv2) (sstatO'$
 $statO sv1 sv2)$
unfolding *match2-12-def validTrans-iff-next* **by** *simp*

lemmas *match2-defsD* = *match2-def match2-1-defD match2-12-defD*

lemma *match12-1-defD*: $match12-1 \Delta s1' s2' statA' sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv1 \wedge \Delta \infty s1' s2' statA' (nextO sv1) sv2 statO$
unfolding *match12-1-def validTrans-iff-next* **by** *simp*

lemma *match12-2-defD*: $match12-2 \Delta s1' s2' statA' sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv2 \wedge \Delta \infty s1' s2' statA' sv1 (nextO sv2) statO$
unfolding *match12-2-def validTrans-iff-next* **by** *simp*

lemma *match12-12-defD*: $match12-12 \Delta s1' s2' statA' sv1 sv2 statO \longleftrightarrow$
 $(let statO' = sstatO' statO sv1 sv2 in$
 $\neg finalV sv1 \wedge \neg finalV sv2 \wedge$
 $(statA' = Diff \longrightarrow statO' = Diff) \wedge$
 $\Delta \infty s1' s2' statA' (nextO sv1) (nextO sv2) statO')$
unfolding *match12-12-def validTrans-iff-next* **by** *simp*

lemmas *match12-defsD* = *match12-def match12-1-defD match12-2-defD match12-12-defD*

lemmas *match-deep-defsD* = *match1-defsD match2-defsD match12-defsD*

lemma *move-1-defD*: $move-1 \Delta w s1 s2 statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv1 \wedge \Delta w s1 s2 statA (nextO sv1) sv2 statO$
unfolding *move-1-def validTrans-iff-next* **by** *simp*

lemma *move-2-defD*: $move-2 \Delta w s1 s2 statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv2 \wedge \Delta w s1 s2 statA sv1 (nextO sv2) statO$
unfolding *move-2-def validTrans-iff-next* **by** *simp*

lemma *move-12-defD*: $move-12 \Delta w s1 s2 statA sv1 sv2 statO \longleftrightarrow$
 $(let statO' = sstatO' statO sv1 sv2 in$
 $\neg finalV sv1 \wedge \neg finalV sv2 \wedge$
 $\Delta w s1 s2 statA (nextO sv1) (nextO sv2) statO')$
unfolding *move-12-def validTrans-iff-next* **by** *simp*

```
lemmas proact-defsD = proact-def move-1-defD move-2-defD move-12-defD
```

```
end
```

```
end
```

4 Unwinding Proof Method for Relative Security

This theory formalizes the notion of unwinding for relative security, and proves its soundness.

```
theory Unwinding
imports Relative-Security
begin
```

4.1 The types and operators underlying unwinding: status, matching operators, etc.

```
context Rel-Sec
begin
```

```
datatype status = Eq | Diff
```

```
fun updStat :: status  $\Rightarrow$  bool  $\times$  'a  $\Rightarrow$  bool  $\times$  'a  $\Rightarrow$  status where
  updStat Eq (True,a) (True,a') = (if a = a' then Eq else Diff)
| updStat stat - - = stat
```

```
definition sstatO' statO sv1 sv2 = updStat statO (isIntV sv1, getObsV sv1)
(isIntV sv2, getObsV sv2)
```

```
definition sstatA' statA s1 s2 = updStat statA (isIntO s1, getObsO s1) (isIntO
s2, getObsO s2)
```

```
lemma updStat-EqI:
```

```
  assumes  $\langle R = S \rangle$ 
```

```
  shows  $\langle \text{updStat Eq } (P, R) (Q, S) = \text{Eq} \rangle$ 
```

```
  apply (cases P)
```

```
  apply (metis assms updStat.simps(1) updStat.simps(4))
```

```
  by (cases Q) auto
```

```
lemma updStat-diff: updStat stat r r = Diff  $\Longrightarrow$  stat = Diff
```

```
  by (metis updStat.elims updStat.simps(1))
```

```
definition initCond ::
```


$(enat \Rightarrow enat \Rightarrow enat \Rightarrow 'stateO \Rightarrow 'stateO \Rightarrow status \Rightarrow 'stateV \Rightarrow 'stateV \Rightarrow status \Rightarrow bool) \Rightarrow bool$ **where**
 $initCond \Delta \equiv \forall s1 s2.$
 $istateO s1 \wedge istateO s2$
 \longrightarrow
 $(\exists sv1 sv2. istateV sv1 \wedge istateV sv2 \wedge corrState sv1 s1 \wedge corrState sv2 s2$
 $\wedge \Delta \infty \infty \infty s1 s2 Eq sv1 sv2 Eq)$

definition $match1-1 \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \equiv$
 $\exists sv1'. validTransV (sv1, sv1') \wedge$
 $\Delta \infty w1 w2 s1' s2 statA sv1' sv2 statO$

definition $match1-12 \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \equiv$
 $(\exists sv1' sv2'.$
 $let statO' = sstatO' statO sv1 sv2 in$
 $validTransV (sv1, sv1') \wedge$
 $validTransV (sv2, sv2') \wedge$
 $\Delta \infty w1 w2 s1' s2 statA sv1' sv2' statO')$

definition $match1 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \equiv$
 $\neg isIntO s1 \longrightarrow$
 $(\forall s1'. validTransO (s1, s1')$
 \longrightarrow
 $(\exists w1' < w1. \exists w2' < w2. \neg isSecO s1 \wedge \Delta \infty w1' w2' s1' s2 statA sv1 sv2$
 $statO) \vee$
 $(\exists w2' < w2. eqSec sv1 s1 \wedge \neg isIntV sv1 \wedge match1-1 \Delta \infty w2' s1 s1' s2$
 $statA sv1 sv2 statO) \vee$
 $(eqSec sv1 s1 \wedge \neg isSecV sv2 \wedge Van.eqAct sv1 sv2 \wedge match1-12 \Delta \infty \infty s1$
 $s1' s2 statA sv1 sv2 statO))$

lemmas $match1-defs = match1-def match1-1-def match1-12-def$

lemma $match1-1-mono:$
 $\Delta \leq \Delta' \Longrightarrow match1-1 \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \Longrightarrow$
 $match1-1 \Delta' w1 w2 s1 s1' s2 statA sv1 sv2 statO$
unfolding $le-fun-def match1-1-def$ **by** $auto$

lemma $match1-12-mono:$
 $\Delta \leq \Delta' \Longrightarrow match1-12 \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \Longrightarrow$
 $match1-12 \Delta' w1 w2 s1 s1' s2 statA sv1 sv2 statO$
unfolding $le-fun-def match1-12-def$ **by** $fastforce$

lemma $match1-mono:$
assumes $\Delta \leq \Delta'$
shows $match1 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \Longrightarrow match1 \Delta' w1 w2 s1 s2$
 $statA sv1 sv2 statO$

unfolding *match1-def* **apply** *clarify subgoal for s1'* **apply**(*erule alle[of - s1']*)
using *match1-1-mono*[*OF assms, of - - s1 s1' s2 statA sv1 sv2 statO*]
match1-12-mono[*OF assms, of - - s1 s1' s2 statA sv1 sv2 statO*]
assms[*unfolded le-fun-def, rule-format, of - - s1' s2 statA sv1 sv2 statO*]
by *fastforce* .

definition *match2-1* $\Delta w1 w2 s1 s2 s2' statA sv1 sv2 statO \equiv$
 $\exists sv2'. \text{validTransV } (sv2, sv2') \wedge$
 $\Delta \infty w1 w2 s1 s2' statA sv1 sv2' statO$

definition *match2-12* $\Delta w1 w2 s1 s2 s2' statA sv1 sv2 statO \equiv$
 $\exists sv1' sv2'.$
let *statO' = sstatO' statO sv1 sv2 in*
 $\text{validTransV } (sv1, sv1') \wedge$
 $\text{validTransV } (sv2, sv2') \wedge$
 $\Delta \infty w1 w2 s1 s2' statA sv1' sv2' statO'$

definition *match2* $\Delta w1 w2 s1 s2 statA sv1 sv2 statO \equiv$
 $\neg \text{isIntO } s2 \longrightarrow$
 $(\forall s2'. \text{validTransO } (s2, s2'))$
 \longrightarrow
 $(\exists w1' < w1. \exists w2' < w2. \neg \text{isSecO } s2 \wedge \Delta \infty w1' w2' s1 s2' statA sv1 sv2$
statO) \vee
 $(\exists w1' < w1. \text{eqSec } sv2 s2 \wedge \neg \text{isIntV } sv2 \wedge \text{match2-1 } \Delta w1' \infty s1 s2 s2' statA$
sv1 sv2 statO) \vee
 $(\neg \text{isSecV } sv1 \wedge \text{eqSec } sv2 s2 \wedge \text{Van.eqAct } sv1 sv2 \wedge \text{match2-12 } \Delta \infty \infty s1$
s2 s2' statA sv1 sv2 statO))

lemmas *match2-defs* = *match2-def match2-1-def match2-12-def*

lemma *match2-1-mono*:

$\Delta \leq \Delta' \implies \text{match2-1 } \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \implies \text{match2-1 } \Delta'$
 $w1 w2 s1 s1' s2 statA sv1 sv2 statO$

unfolding *le-fun-def match2-1-def* **by** *auto*

lemma *match2-12-mono*:

$\Delta \leq \Delta' \implies \text{match2-12 } \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \implies \text{match2-12 } \Delta'$
 $w1 w2 s1 s1' s2 statA sv1 sv2 statO$

unfolding *le-fun-def match2-12-def* **by** *fastforce*

lemma *match2-mono*:

assumes $\Delta \leq \Delta'$

shows *match2* $\Delta w1 w2 s1 s2 statA sv1 sv2 statO \implies \text{match2 } \Delta' w1 w2 s1 s2$
statA sv1 sv2 statO

unfolding *match2-def* **apply** *clarify subgoal for s2'* **apply**(*erule alle[of - s2']*)

using *match2-1-mono*[*OF assms, of - - s1 s2 s2' statA sv1 sv2 statO*]

match2-12-mono[*OF assms, of - - s1 s2 s2' statA sv1 sv2 statO*]

assms[*unfolded le-fun-def, rule-format, of - - - s1 s2' statA sv1 sv2 statO*]
by *fastforce* .

definition *match12-1* $\Delta w1 w2 s1' s2' statA' sv1 sv2 statO \equiv$
 $\exists sv1'. \text{validTransV } (sv1, sv1') \wedge$
 $\Delta \infty w1 w2 s1' s2' statA' sv1' sv2 statO$

definition *match12-2* $\Delta w1 w2 s1' s2' statA' sv1 sv2 statO \equiv$
 $\exists sv2'. \text{validTransV } (sv2, sv2') \wedge$
 $\Delta \infty w1 w2 s1' s2' statA' sv1 sv2' statO$

definition *match12-12* $\Delta w1 w2 s1' s2' statA' sv1 sv2 statO \equiv$
 $\exists sv1' sv2'.$
let *statO' = sstatO' statO sv1 sv2 in*
 $\text{validTransV } (sv1, sv1') \wedge$
 $\text{validTransV } (sv2, sv2') \wedge$
 $(statA' = \text{Diff} \longrightarrow statO' = \text{Diff}) \wedge$
 $\Delta \infty w1 w2 s1' s2' statA' sv1' sv2' statO'$

definition *match12* $\Delta w1 w2 s1 s2 statA sv1 sv2 statO \equiv$
 $\forall s1' s2'.$
let *statA' = sstatA' statA s1 s2 in*
 $\text{validTransO } (s1, s1') \wedge$
 $\text{validTransO } (s2, s2') \wedge$
 $\text{Opt.eqAct } s1 s2 \wedge$
 $\text{isIntO } s1 \wedge \text{isIntO } s2$
 \longrightarrow
 $(\exists w1' < w1. \exists w2' < w2. \neg \text{isSecO } s1 \wedge \neg \text{isSecO } s2 \wedge (statA = statA' \vee statO$
 $= \text{Diff}) \wedge$
 $\Delta \infty w1' w2' s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(\exists w2' < w2. \neg \text{isSecO } s2 \wedge \text{eqSec } sv1 s1 \wedge \neg \text{isIntV } sv1 \wedge$
 $(statA = statA' \vee statO = \text{Diff}) \wedge$
 $\text{match12-1 } \Delta \infty w2' s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(\exists w1' < w1. \neg \text{isSecO } s1 \wedge \text{eqSec } sv2 s2 \wedge \neg \text{isIntV } sv2 \wedge$
 $(statA = statA' \vee statO = \text{Diff}) \wedge$
 $\text{match12-2 } \Delta w1' \infty s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(\text{eqSec } sv1 s1 \wedge \text{eqSec } sv2 s2 \wedge \text{Van.eqAct } sv1 sv2 \wedge$
 $\text{match12-12 } \Delta \infty \infty s1' s2' statA' sv1 sv2 statO)$

lemmas *match12-defs* = *match12-def match12-1-def match12-2-def match12-12-def*

lemma *match12-simpleI*:
assumes $\bigwedge s1' s2' statA'.$

$statA' = sstatA' statA s1 s2 \implies$
 $validTransO (s1, s1') \implies$
 $validTransO (s2, s2') \implies$
 $Opt.eqAct s1 s2 \implies$
 $(\exists w1' < w1. \exists w2' < w2. \neg isSecO s1 \wedge \neg isSecO s2 \wedge (statA = statA' \vee statO$
 $= Diff) \wedge$
 $\Delta \infty w1' w2' s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(eqSec sv1 s1 \wedge eqSec sv2 s2 \wedge Van.eqAct sv1 sv2 \wedge$
 $match12-12 \Delta \infty \infty s1' s2' statA' sv1 sv2 statO)$
shows $match12 \Delta w1 w2 s1 s2 statA sv1 sv2 statO$
using *assms unfolding match12-def Let-def by blast*

lemma *match12-1-mono*:
 $\Delta \leq \Delta' \implies match12-1 \Delta w1 w2 s1' s2' statA' sv1 sv2 statO \implies match12-1 \Delta'$
 $w1 w2 s1' s2' statA' sv1 sv2 statO$
unfolding *le-fun-def match12-1-def by auto*

lemma *match12-2-mono*:
 $\Delta \leq \Delta' \implies match12-2 \Delta w1 w2 s1 s2' statA' sv1 sv2 statO \implies match12-2 \Delta'$
 $w1 w2 s1 s2' statA' sv1 sv2 statO$
unfolding *le-fun-def match12-2-def by auto*

lemma *match12-12-mono*:
 $\Delta \leq \Delta' \implies match12-12 \Delta w1 w2 s1' s2' statA' sv1 sv2 statO \implies match12-12$
 $\Delta' w1 w2 s1' s2' statA' sv1 sv2 statO$
unfolding *le-fun-def match12-12-def by fastforce*

lemma *match12-mono*:
assumes $\Delta \leq \Delta'$
shows $match12 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \implies match12 \Delta' w1 w2 s1 s2$
 $statA sv1 sv2 statO$
unfolding *match12-def apply clarify subgoal for s1' s2' apply(erule allE[of -*
 $s1']) apply(erule allE[of - s2'])$
using *match12-1-mono[OF assms, of - - s1' s2' - sv1 sv2 statO]*
 $match12-2-mono[OF assms, of - - s1' s2' - sv1 sv2 statO]$
 $match12-12-mono[OF assms, of - - s1' s2' - sv1 sv2 statO]$
 $assms[unfolding le-fun-def, rule-format, of - - - s1' s2'$
 $sstatA' statA s1 s2 sv1 sv2 statO]$
apply *simp by blast .*

definition *match* $\Delta w1 w2 s1 s2 statA sv1 sv2 statO \equiv$
 $match1 \Delta w1 w2 s1 s2 statA sv1 sv2 statO$
 \wedge
 $match2 \Delta w1 w2 s1 s2 statA sv1 sv2 statO$
 \wedge
 $match12 \Delta w1 w2 s1 s2 statA sv1 sv2 statO$

lemmas *match-defs* = *match1-def match2-def match12-def*
lemmas *match-deep-defs* = *match1-defs match2-defs match12-defs*

lemma *match-mono*:

assumes $\Delta \leq \Delta'$

shows *match* Δ *w1 w2 s1 s2 statA sv1 sv2 statO* \implies *match* Δ' *w1 w2 s1 s2 statA sv1 sv2 statO*

unfolding *match-def* **using** *match1-mono*[*OF assms*] *match2-mono*[*OF assms*] *match12-mono*[*OF assms*] **by** *auto*

definition *move-1* Δ *w w1 w2 s1 s2 statA sv1 sv2 statO* \equiv

$\exists sv1'. \text{validTransV } (sv1, sv1') \wedge$
 Δ *w w1 w2 s1 s2 statA sv1' sv2 statO*

definition *move-2* Δ *w w1 w2 s1 s2 statA sv1 sv2 statO* \equiv

$\exists sv2'. \text{validTransV } (sv2, sv2') \wedge$
 Δ *w w1 w2 s1 s2 statA sv1 sv2' statO*

definition *move-12* Δ *w w1 w2 s1 s2 statA sv1 sv2 statO* \equiv

$\exists sv1' sv2'.$
let *statO'* = *sstatO' statO sv1 sv2 in*
 $\text{validTransV } (sv1, sv1') \wedge \text{validTransV } (sv2, sv2') \wedge$
 Δ *w w1 w2 s1 s2 statA sv1' sv2' statO'*

definition *proact* Δ *w w1 w2 s1 s2 statA sv1 sv2 statO* \equiv

$(\neg \text{isSecV } sv1 \wedge \neg \text{isIntV } sv1 \wedge \text{move-1 } \Delta$ *w w1 w2 s1 s2 statA sv1 sv2 statO*)
 \vee
 $(\neg \text{isSecV } sv2 \wedge \neg \text{isIntV } sv2 \wedge \text{move-2 } \Delta$ *w w1 w2 s1 s2 statA sv1 sv2 statO*)
 \vee
 $(\neg \text{isSecV } sv1 \wedge \neg \text{isSecV } sv2 \wedge \text{Van.eqAct } sv1 sv2 \wedge \text{move-12 } \Delta$ *w w1 w2 s1 s2 statA sv1 sv2 statO*)

lemmas *proact-defs* = *proact-def move-1-def move-2-def move-12-def*

lemma *move-1-mono*:

$\Delta \leq \Delta' \implies \text{move-1 } \Delta$ *w w1 w2 s1 s2 statA sv1 sv2 statO* $\implies \text{move-1 } \Delta'$ *w w1 w2 s1 s2 statA sv1 sv2 statO*

unfolding *le-fun-def move-1-def* **by** *auto*

lemma *move-2-mono*:

$\Delta \leq \Delta' \implies \text{move-2 } \Delta$ *w w1 w2 s1 s2 statA sv1 sv2 statO* $\implies \text{move-2 } \Delta'$ *w w1 w2 s1 s2 statA sv1 sv2 statO*

unfolding *le-fun-def move-2-def* **by** *auto*

lemma *move-12-mono*:

$\Delta \leq \Delta' \implies \text{move-12 } \Delta$ *w w1 w2 s1 s2 statA sv1 sv2 statO* $\implies \text{move-12 } \Delta'$ *w w1*

$w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
unfolding *le-fun-def move-12-def* **by** *fastforce*

lemma *proact-mono*:

assumes $\Delta \leq \Delta'$

shows $proact\ \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies proact\ \Delta'\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

unfolding *proact-def* **using** *move-1-mono[OF assms]* *move-2-mono[OF assms]* *move-12-mono[OF assms]* **by** *auto*

4.2 The definition of unwinding

definition *unwindCond* ::

$(enat \Rightarrow enat \Rightarrow enat \Rightarrow 'stateO \Rightarrow 'stateO \Rightarrow status \Rightarrow 'stateV \Rightarrow 'stateV \Rightarrow status \Rightarrow bool) \Rightarrow bool$

where

$unwindCond\ \Delta \equiv \forall w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$

$reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$
 $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

\longrightarrow

$(finalO\ s1 \longleftrightarrow finalO\ s2) \wedge (finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO\ s2)$

\wedge

$(statA = Eq \longrightarrow (isIntO\ s1 \longleftrightarrow isIntO\ s2))$

\wedge

$((\exists v < w. proact\ \Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO)$

\vee

$match\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

$)$

lemma *unwindCond-simpleI*:

assumes

$\bigwedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$

$reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$

$\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

\implies

$(finalO\ s1 \longleftrightarrow finalO\ s2) \wedge (finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO\ s2)$

and

$\bigwedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$

$reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$

$\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies statA = Eq$

\implies

$isIntO\ s1 \longleftrightarrow isIntO\ s2$

and

$\bigwedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$

$reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
 \implies
 $match\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
shows *unwindCond* Δ
using *assms* **unfolding** *unwindCond-def* **by** *auto*

4.3 The soundness of unwinding

The proof of soundness for general unwinding is significantly more elaborate than that for the finitary case.

definition $\psi\ s1\ tr1\ s2\ tr2\ statO\ sv1\ trv1\ sv2\ trv2 \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge$
 $Van.validFromS\ sv2\ trv2 \wedge$
 $(finalV\ (lastt\ sv1\ trv1) \longleftrightarrow finalO\ (lastt\ s1\ tr1)) \wedge (finalV\ (lastt\ sv2\ trv2) \longleftrightarrow$
 $finalO\ (lastt\ s2\ tr2)) \wedge$
 $Van.S\ trv1 = Opt.S\ tr1 \wedge Van.S\ trv2 = Opt.S\ tr2 \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $(statO = Eq \wedge Opt.O\ tr1 \neq Opt.O\ tr2 \longrightarrow Van.O\ trv1 \neq Van.O\ trv2)$

lemma ψ -*completedFrom*: $completedFromO\ s1\ tr1 \implies completedFromO\ s2\ tr2 \implies$

$\psi\ s1\ tr1\ s2\ tr2\ statO\ sv1\ trv1\ sv2\ trv2$
 $\implies completedFromV\ sv1\ trv1 \wedge completedFromV\ sv2\ trv2$

unfolding ψ -*def* *Opt.completedFrom-def* *Van.completedFrom-def* *lastt-def*
by *presburger*

lemma *completedFromO-lastt*: $completedFromO\ s1\ tr1 \implies finalO\ (lastt\ s1\ tr1)$

unfolding *Opt.completedFrom-def* *lastt-def* **by** *auto*

lemma *rsecure-strong*:

assumes

$\bigwedge s1\ tr1\ s2\ tr2.$

$istateO\ s1 \wedge Opt.validFromS\ s1\ tr1 \wedge completedFromO\ s1\ tr1 \wedge$

$istateO\ s2 \wedge Opt.validFromS\ s2\ tr2 \wedge completedFromO\ s2\ tr2 \wedge$

$Opt.A\ tr1 = Opt.A\ tr2$

\implies

$\exists sv1\ trv1\ sv2\ trv2.$

$istateV\ sv1 \wedge istateV\ sv2 \wedge corrState\ sv1\ s1 \wedge corrState\ sv2\ s2 \wedge$

$\psi\ s1\ tr1\ s2\ tr2\ Eq\ sv1\ trv1\ sv2\ trv2$

shows *rsecure*

unfolding *rsecure-def2* **apply** *safe*

subgoal for $s1\ tr1\ s2\ tr2$

using *assms*[*of* $s1\ tr1\ s2\ tr2$]

using ψ -*completedFrom* ψ -*def* *completedFromO-lastt* **apply** *clarsimp* **by** *metis* .

proposition *unwindCond-ex-ψ*:
assumes *unwind*: *unwindCond* Δ
and Δ: Δ *w w1 w2 s1 s2 statA sv1 sv2 statO* **and** *stat*: (*statA* = *Diff* \longrightarrow *statO* = *Diff*)
and *v*: *Opt.validFromS s1 tr1 Opt.completedFrom s1 tr1 Opt.validFromS s2 tr2 Opt.completedFrom s2 tr2*
and *tr14*: *Opt.A tr1 = Opt.A tr2*
and *r*: *reachO s1 reachO s2 reachV sv1 reachV sv2*
shows \exists *trv1 trv2. ψ s1 tr1 s2 tr2 statO sv1 trv1 sv2 trv2*
using *assms(2-)*
proof(*induction length tr1 + length tr2 w*
arbitrary: w1 w2 s1 s2 statA sv1 sv2 statO tr1 tr2 rule: less2-induct')
case (*less w tr1 tr2 w1 w2 s1 s2 statA sv1 sv2 statO*)
note *ok* = \langle *statA = Diff* \longrightarrow *statO = Diff* \rangle
note Δ = \langle Δ *w w1 w2 s1 s2 statA sv1 sv2 statO* \rangle
note *A34* = \langle *Opt.A tr1 = Opt.A tr2* \rangle
note *r34* = *less.prem(8,9)* **note** *r12* = *less.prem(10,11)*
note *r* = *r34 r12*
note *r3* = *r34(1)* **note** *r4* = *r34(2)* **note** *r1* = *r12(1)* **note** *r2* = *r12(2)*

have *i34*: *statA = Eq* \longrightarrow *isIntO s1 = isIntO s2*
and *f34*: *finalO s1 = finalO s2* \wedge *finalV sv1 = finalO s1* \wedge *finalV sv2 = finalO s2*
using Δ *unwind[unfolded unwindCond-def]* *r* **by** *auto*

have *proact-match*: (\exists *v* < *w*. *proact* Δ *v w1 w2 s1 s2 statA sv1 sv2 statO*) \vee *match* Δ *w1 w2 s1 s2 statA sv1 sv2 statO*
using Δ *unwind[unfolded unwindCond-def]* *r* **by** *auto*
show ?*case* **using** *proact-match* **proof** *safe*
fix *v* **assume** *v*: *v* < *w*
assume *proact* Δ *v w1 w2 s1 s2 statA sv1 sv2 statO*
thus ?*thesis* **unfolding** *proact-def* **proof** *safe*
assume *sv1*: \neg *isSecV sv1* \neg *isIntV sv1* **and** *move-1* Δ *v w1 w2 s1 s2 statA sv1 sv2 statO*
then obtain *sv1'*
where *0*: *validTransV (sv1,sv1')*
and Δ: Δ *v w1 w2 s1 s2 statA sv1' sv2 statO*
unfolding *move-1-def* **by** *auto*
have *r1'*: *reachV sv1'* **using** *r1 0* **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain *trv1 trv2* **where** *ψ*: *ψ s1 tr1 s2 tr2 statO sv1' trv1 sv2 trv2*
using *less(2)[OF v, of tr1 tr2 w1 w2 s1 s2 statA sv1' sv2 statO, simplified, OF Δ ok - - - - r34 r1' r2]*
using *A34 less.prem(3-6)* **by** *blast*
show ?*thesis* **apply**(*rule exI[of - sv1 # trv1]*) **apply**(*rule exI[of - trv2]*)
using *ψ ok 0 sv1 unfolding ψ-def Van.completedFrom-def* **by** *auto*
next
assume *sv2*: \neg *isSecV sv2* \neg *isIntV sv2* **and** *move-2* Δ *v w1 w2 s1 s2 statA sv1 sv2 statO*


```

then obtain sv2'
where 0: validTransV (sv2,sv2')
and Δ: Δ v w1 w2 s1 s2 statA sv1 sv2' statO
unfolding move-2-def by auto
have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
obtain trv1 trv2 where ψ: ψ s1 tr1 s2 tr2 statO sv1 trv1 sv2' trv2
using less(2)[OF v, of tr1 tr2 w1 w2 s1 s2 statA sv1 sv2' statO, simplified,
OF Δ ok - - - - r34 r1 r2']
using A34 less.prem(3-6) by blast
show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using ψ ok 0 sv2 unfolding ψ-def Van.completedFrom-def by auto
next
assume sv12: ¬ isSecV sv1 ¬ isSecV sv2 Van.eqAct sv1 sv2
and move-12 Δ v w1 w2 s1 s2 statA sv1 sv2 statO
then obtain sv1' sv2' statO'
where 0: statO' = sstatO' statO sv1 sv2
validTransV (sv1,sv1') ¬ isSecV sv1
validTransV (sv2,sv2') ¬ isSecV sv2
Van.eqAct sv1 sv2
and Δ: Δ v w1 w2 s1 s2 statA sv1' sv2' statO'
unfolding move-12-def by auto
have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
fst-conv snd-conv)+
have ok': statA = Diff → statO' = Diff using ok 0 unfolding sstatO'-def
by (cases statO, auto)
obtain trv1 trv2 where ψ: ψ s1 tr1 s2 tr2 statO' sv1' trv1 sv2' trv2
using less(2)[OF v, of tr1 tr2 w1 w2 s1 s2 statA sv1' sv2' statO', simplified,
OF Δ ok' - - - - r34 r12']
using A34 less.prem(3-6) by blast
show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
using ψ ok' 0 sv12 unfolding ψ-def sstatO'-def Van.completedFrom-def
using Van.A.Cons-unfold Van.eqAct-def completedFromO-lastt less.prem(4)
less.prem(6) by auto
qed
next
assume m: match Δ w1 w2 s1 s2 statA sv1 sv2 statO
show ?thesis
proof(cases length tr1 ≤ Suc 0)
case True note tr1 = True
hence tr1 = [] ∨ tr1 = [s1]
by (metis Simple-Transition-System.validFromS-Cons-iff Suc-length-conv le-Suc-eq
le-zero-eq length-0-conv less.prem(3))
hence finalO s1 using less(3-6)
using Opt.completed-Cons Opt.completed-Nil by blast
hence f4: finalO s2 using f34 by blast
hence tr2: tr2 = [] ∨ tr2 = [s2]
by (metis Opt.final-def Simple-Transition-System.validFromS-Cons-iff less.prem(5)
neq-Nil-conv)

```

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show ?thesis apply(rule exI[of - [sv1]], rule exI[of - [sv2]]) using tr1 tr2
using f4 f34
using completedFromO-lastt less.premis(4)
by (auto simp add: lastt-def  $\psi$ -def)
next
case False
then obtain s13 tr1' where tr1: tr1 = s13 # tr1' and tr1'NE: tr1'  $\neq$  []
  by (cases tr1, auto)
have s13[simp]: s13 = s1 using  $\langle$ Opt.validFromS s1 tr1 $\rangle$ 
  by (simp add: Opt.validFromS-Cons-iff tr1)
obtain s1' where
  trn3: validTransO (s1,s1') and
  tr1': Opt.validFromS s1' tr1' using  $\langle$ Opt.validFromS s1 tr1 $\rangle$ 
unfolding tr1 s13 by (metis tr1'NE Simple-Transition-System.validFromS-Cons-iff)
  have r3': reachO s1' using r3 trn3 by (metis Opt.reach.Step fst-conv
snd-conv)
  have f3:  $\neg$  finalO s1 using Opt.final-def trn3 by blast
  hence f4:  $\neg$  finalO s2 using f34 by blast
  hence tr2:  $\neg$  length tr2  $\leq$  Suc 0
  by (metis Opt.completed-Cons Simple-Transition-System.validFromS-Cons-iff

bot-nat-0.extremum completedFromO-def length-Cons less.premis(5) less.premis(6)
neq-Nil-conv not-less-eq-eq)

then obtain s24 tr2' where tr2: tr2 = s24 # tr2' and tr2'NE: tr2'  $\neq$  []
by (cases tr2, auto)
have s24[simp]: s24 = s2 using  $\langle$ Opt.validFromS s2 tr2 $\rangle$ 
by (simp add: Opt.validFromS-Cons-iff tr2)
obtain s2' where
  trn4: validTransO (s2,s2')  $\vee$  (s2 = s2'  $\wedge$  tr2' = []) and
  tr2': Opt.validFromS s2' tr2' using  $\langle$ Opt.validFromS s2 tr2 $\rangle$ 
unfolding tr2 s24 using Opt.validFromS-Cons-iff by auto
have r34': reachO s1' reachO s2'
using r3 trn3 r4 trn4 by (metis Opt.reach.Step fst-conv snd-conv)+
note r3' = r34'(1) note r4' = r34'(2)
define statA' where statA': statA' = sstatA' statA s1 s2
have  $\neg$  isIntO s1  $\vee$   $\neg$  isIntO s2  $\vee$  (isIntO s1  $\wedge$  isIntO s2)
by auto
thus ?thesis
proof safe
  assume isAO3:  $\neg$  isIntO s1
  have O33': Opt.O tr1 = Opt.O tr1' Opt.A tr1 = Opt.A tr1'
  using isAO3 unfolding tr1 by auto
  have A34': Opt.A tr1' = Opt.A tr2
  using A34 O33'(2) by auto
  have m: match1  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
match-def by auto
  have ( $\exists$  w1' < w1.  $\exists$  w2' < w2.  $\neg$  isSecO s1  $\wedge$   $\Delta$   $\infty$  w1' w2' s1' s2 statA sv1
sv2 statO)  $\vee$ 

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      (∃ w2' < w2. eqSec sv1 s1 ∧ ¬ isIntV sv1 ∧ match1-1 Δ ∞ w2' s1 s1'
s2 statA sv1 sv2 statO) ∨
      (eqSec sv1 s1 ∧ ¬ isSecV sv2 ∧ Van.eqAct sv1 sv2 ∧ match1-12 Δ ∞
∞ s1 s1' s2 statA sv1 sv2 statO)
using m isAO3 trn3 ok unfolding match1-def by auto
thus ?thesis
proof safe
  fix w1' w2'
  assume ¬ isSecO s1 and Δ: Δ ∞ w1' w2' s1' s2 statA sv1 sv2 statO
  hence S3: Opt.S tr1' = Opt.S tr1 unfolding tr1 by auto
  obtain trv1 trv2 where ψ: ψ s1 tr1' s2 tr2 statO sv1 trv1 sv2 trv2
    using less(1)[of tr1' tr2, OF - Δ - - - - - r3' r4 r12, unfolded O33',
simplified]
  using less.prem1 tr1' ok A34' f3 f4 unfolding tr1 Opt.completedFrom-def
by (auto split: if-splits simp: ψ-def lastt-def)
  show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
  using ψ O33' S3 unfolding ψ-def
  using completedFromO-lastt less.prem1(4)
by (auto simp add: tr1 tr1'NE)
next
  fix w2'
  assume trn13: eqSec sv1 s1 and
Atrn1: ¬ isIntV sv1 and match1-1 Δ ∞ w2' s1 s1' s2 statA sv1 sv2 statO
  then obtain sv1' where
    trn1: validTransV (sv1, sv1') and
    Δ: Δ ∞ ∞ w2' s1' s2 statA sv1' sv2 statO
  unfolding match1-1-def by auto
    have r1': reachV sv1' using r1 trn1 by (metis Van.reach.Step fst-conv
snd-conv)
  obtain trv1 trv2 where ψ: ψ s1 tr1' s2 tr2 statO sv1' trv1 sv2 trv2
    using less(1)[of tr1' tr2, OF - Δ - - - - - r3' r4 r1' r2, unfolded O33',
simplified]
  using less.prem1 tr1' ok A34' f3 f4 unfolding tr1 tr2 Opt.completedFrom-def

by (auto simp: ψ-def lastt-def split: if-splits)
show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
using ψ O33' unfolding tr1 tr2 Van.completedFrom-def
using Van.validFromS-Cons trn1 tr1'NE tr2'NE
using isAO3 ok Atrn1 eqSec-S-Cons trn13
unfolding ψ-def using completedFromO-lastt less.prem1(4) tr1 by auto

next
assume sv2: ¬ isSecV sv2 and trn13: eqSec sv1 s1 and
Atrn12: Van.eqAct sv1 sv2 and match1-12 Δ ∞ ∞ s1 s1' s2 statA sv1
sv2 statO
then obtain sv1' sv2' statO' where
  statO': statO' = sstatO' statO sv1 sv2 and
  trn1: validTransV (sv1, sv1') and
  trn2: validTransV (sv2, sv2') and

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    Δ: Δ ∞ ∞ ∞ s1' s2 statA sv1' sv2' statO'
    unfolding match1-12-def by auto
    have r12': reachV sv1' reachV sv2'
    using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
    obtain trv1 trv2 where ψ: ψ s1' tr1' s2 tr2 statO' sv1' trv1 sv2' trv2
      using less(1)[of tr1' tr2, OF - Δ - - - - - r3' r4' r12', unfolded O33',
simplified]
    using less.premis tr1' ok A34' f3 f4 unfolding tr1 tr2 Opt.completedFrom-def
statO' sstatO'-def
    by auto presburger+
    show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
      using ψ O33' tr1'NE tr2'NE sv2
      using Van.validFromS-Cons trn1 trn2
      using isAO3 ok Atrn12 eqSec-S-Cons trn13 f3 f34 s13
    unfolding ψ-def tr1 Van.completedFrom-def Van.eqAct-def statO' sstatO'-def
      using Van.A.Cons-unfold tr1' trn3 by auto
    qed
  next
  assume isAO4: ¬ isIntO s2
  have O44': Opt.O tr2 = Opt.O tr2' Opt.A tr2 = Opt.A tr2'
  using isAO4 unfolding tr2 by auto
  have A34': Opt.A tr1 = Opt.A tr2'
  using A34 O44'(2) by auto
  have m: match2 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
match-def by auto
  have (∃ w1' < w1. ∃ w2' < w2. ¬ isSecO s2 ∧ Δ ∞ w1' w2' s1 s2' statA sv1
sv2 statO) ∨
    (∃ w1' < w1. eqSec sv2 s2 ∧ ¬ isIntV sv2 ∧ match2-1 Δ w1' ∞ s1 s2
s2' statA sv1 sv2 statO) ∨
    (¬ isSecV sv1 ∧ eqSec sv2 s2 ∧ Van.eqAct sv1 sv2 ∧ match2-12 Δ ∞
∞ s1 s2 s2' statA sv1 sv2 statO)
  using m isAO4 trn4 ok tr2'NE unfolding match2-def by auto
  thus ?thesis
  proof safe
    fix w1' w2'
    assume ¬ isSecO s2 and Δ: Δ ∞ w1' w2' s1 s2' statA sv1 sv2 statO
    hence S4: Opt.S tr2' = Opt.S tr2 unfolding tr2 by auto
    obtain trv1 trv2 where ψ: ψ s1 tr1 s2' tr2' statO sv1 trv1 sv2 trv2
      using less(1)[of tr1 tr2', OF - Δ - - - - - r3' r4', simplified]
    using less.premis tr2' ok A34' tr1'NE tr2'NE unfolding tr1 tr2 Opt.completedFrom-def
by (cases isIntO s2, auto)
    show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
      using ψ O44' S4 unfolding ψ-def
      using completedFromO-lastt less.premis(6)
      unfolding Opt.completedFrom-def using tr2 tr2'NE by auto
  next
  fix w1'
  assume trn24: eqSec sv2 s2 and

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Atrn2: \neg isIntV sv2 **and** match2-1 $\Delta w1' \infty s1 s2 s2' statA sv1 sv2 statO$
then obtain sv2' **where** trn2: validTransV (sv2,sv2') **and**
 Δ : $\Delta \infty w1' \infty s1 s2' statA sv1 sv2' statO$
unfolding match2-1-def **by** auto
have r2': reachV sv2' **using** r2 trn2 **by** (metis Van.reach.Step fst-conv
snd-conv)
obtain trv1 trv2 **where** ψ : $\psi s1 tr1 s2' tr2' statO sv1 trv1 sv2' trv2$
using less(1)[of tr1 tr2', OF - Δ - - - - - r3 r4' r1 r2', simplified]
using less.premis tr2' ok A34' tr1'NE tr2'NE **unfolding** tr1 tr2 Opt.completedFrom-def
by (cases isIntO s2, auto)
show ?thesis **apply**(rule exI[of - trv1]) **apply**(rule exI[of - sv2 # trv2])
using ψ tr1'NE tr2'NE
using Van.validFromS-Cons trn2
using isAO4 ok Atrn2 eqSec-S-Cons trn24
unfolding ψ -def tr1 tr2 s13 s24 Van.completedFrom-def lastt-def **by** auto
next
assume sv1: \neg isSecV sv1 **and** trn24: eqSec sv2 s2 **and**
Atrn12: Van.eqAct sv1 sv2 **and** match2-12 $\Delta \infty \infty s1 s2 s2' statA sv1$
sv2 statO
then obtain sv1' sv2' statO' **where**
statO': statO' = sstatO' statO sv1 sv2 **and**
trn1: validTransV (sv1,sv1') **and**
trn2: validTransV (sv2,sv2') **and**
 Δ : $\Delta \infty \infty \infty s1 s2' statA sv1' sv2' statO'$
unfolding match2-12-def **by** auto
have r12': reachV sv1' reachV sv2'
using r1 trn1 r2 trn2 **by** (metis Van.reach.Step fst-conv snd-conv)+
obtain trv1 trv2 **where** ψ : $\psi s1 tr1 s2' tr2' statO' sv1' trv1 sv2' trv2$
using less(1)[of tr1 tr2', OF - Δ - - - - - r3 r4' r12', simplified]
using less.premis tr2' ok A34' tr1'NE tr2'NE **unfolding** tr1 tr2 Opt.completedFrom-def
statO' sstatO'-def
by (cases isIntO s2, auto)
show ?thesis **apply**(rule exI[of - sv1 # trv1]) **apply**(rule exI[of - sv2 #
trv2])
using ψ O44' tr1'NE tr2'NE sv1
using Van.validFromS-Cons trn1 trn2
using isAO4 ok Atrn12 eqSec-S-Cons trn24
unfolding ψ -def tr2 tr1'NE Van.completedFrom-def Van.eqAct-def
statO' sstatO'-def
using Van.A.Cons-unfold tr2' trn4 **by** auto
qed
next
assume isAO34: isIntO s1 isIntO s2
have A34': getActO s1 = getActO s2 Opt.A tr1' = Opt.A tr2'
using A34 isAO34 tr1'NE tr2'NE **unfolding** tr1 tr2 **by** auto
have O33': Opt.O tr1 = getObsO s1 # Opt.O tr1' **and**
O44': Opt.O tr2 = getObsO s2 # Opt.O tr2'
using isAO34 tr1'NE tr2'NE **unfolding** tr1 s13 tr2 s24 **by** auto
have m: match12 $\Delta w1 w2 s1 s2 statA sv1 sv2 statO$ **using** m **unfolding**

statA' match-def by auto
have *trn34*: *getObsO s1 = getObsO s2* \vee *statA' = Diff*
using *isAO34 unfolding statA' sstatA'-def* by (*cases statA, auto*)
have ($\exists w1' < w1. \exists w2' < w2. \neg isSecO s1 \wedge \neg isSecO s2 \wedge (statA = statA'$
 $\vee statO = Diff) \wedge$
 $\Delta \infty w1' w2' s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(\exists w2' < w2. \neg isSecO s2 \wedge eqSec sv1 s1 \wedge \neg isIntV sv1 \wedge$
 $(statA = statA' \vee statO = Diff) \wedge$
 $match12-1 \Delta \infty w2' s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(\exists w1' < w1. \neg isSecO s1 \wedge eqSec sv2 s2 \wedge \neg isIntV sv2 \wedge$
 $(statA = statA' \vee statO = Diff) \wedge$
 $match12-2 \Delta w1' \infty s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(eqSec sv1 s1 \wedge eqSec sv2 s2 \wedge Van.eqAct sv1 sv2 \wedge$
 $match12-12 \Delta \infty \infty s1' s2' statA' sv1 sv2 statO)$
(is ?K1 \vee ?K2 \vee ?K3 \vee ?K4)
using *m[unfolded match12-def, rule-format, of s1' s2']*
isAO34 A34' trn3 trn4 tr1'NE tr2'NE
unfolding *s13 s24 trn34 statA' Opt.eqAct-def sstatA'-def* by auto
thus *?thesis proof (elim disjE)*
assume *K1: ?K1*
then obtain *w1' w2'* where $\Delta: \Delta \infty w1' w2' s1' s2' statA' sv1 sv2 statO$
by auto
have *ok'*: (*statA' = Diff* \longrightarrow *statO = Diff*)
using *ok K1 unfolding statA' using isAO34* by auto
obtain *trv1 trv2* where $\psi: \psi s1' tr1' s2' tr2' statO sv1 trv1 sv2 trv2$
using *less(1)[of tr1' tr2', OF - Δ - - - - - r34' r12, simplified]*
using *less.premis tr1' tr2' ok' A34' isAO34 tr1'NE tr2'NE unfolding tr1*
tr2 Opt.completedFrom-def by auto
show *?thesis apply(rule exI[of - trv1]) apply(rule exI[of - trv2])*
using ψ *trn34 O33' O44' K1 ok unfolding ψ -def tr1 tr2*
using *completedFromO-lastt less.premis(4,6)*
unfolding *Opt.completedFrom-def using tr1 tr2 tr1'NE tr2'NE* by auto
next
assume *K2: ?K2*
then obtain *w2' sv1'* where
trn1: validTransV (sv1, sv1') and
trn13: eqSec sv1 s1 and
Atrn1: $\neg isIntV sv1$ and *ok': (statA' = statA \vee statO = Diff)* and
 $\Delta: \Delta \infty \infty w2' s1' s2' statA' sv1' sv2 statO$
unfolding *match12-1-def* by auto
have *r1'*: *reachV sv1'* **using** *r1 trn1* by (*metis Van.reach.Step fst-conv*
snd-conv)
obtain *trv1 trv2* where $\psi: \psi s1' tr1' s2' tr2' statO sv1' trv1 sv2 trv2$
using *less(1)[of tr1' tr2', OF - Δ - - - - - r34' r1' r2, simplified]*
using *less.premis tr1' tr2' ok' A34' tr1'NE tr2'NE unfolding tr1 tr2*
Opt.completedFrom-def by auto

```

show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
using  $\psi$  O33' O44' tr1'NE tr2'NE unfolding tr1 tr2
using Van.validFromS-Cons trn1 ok
using K2 ok' Atrn1 eqSec-S-Cons trn13 trn34
unfolding statA' Van.completedFrom-def eqSec-def
using s13 tr1 tr1' tr2' trn3 trn4
by simp (smt (verit, best) Opt.S.Cons-unfold Simple-Transition-System.lastt-Cons

Van.A.Cons-unfold Van.O.Cons-unfold  $\psi$ -def completedFromO-lastt f3 f34
lastt-Nil
less.premis(4) status.simps(1))
next
assume K3: ?K3
then obtain w1' sv2' where
trn2: validTransV (sv2,sv2') and
trn24: eqSec sv2 s2 and
Atrn2:  $\neg$  isIntV sv2 and ok': (statA' = statA  $\vee$  statO = Diff) and
 $\Delta$ :  $\Delta \infty w1' \infty s1' s2' statA' sv1 sv2' statO$ 
unfolding match12-2-def by auto
have r2': reachV sv2' using r2 trn2 by (metis Van.reach.Step fst-conv
snd-conv)
obtain trv1 trv2 where  $\psi$ :  $\psi s1' tr1' s2' tr2' statO sv1 trv1 sv2' trv2$ 
using less(1)[of tr1' tr2', OF -  $\Delta$  - - - - - r34' r1 r2', simplified]
using less.premis tr1' tr2' ok' A34' tr1'NE tr2'NE unfolding tr1 tr2
Opt.completedFrom-def by auto
show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using  $\psi$  O33' O44' tr1'NE tr2'NE unfolding  $\psi$ -def tr1 tr2
using Van.validFromS-Cons trn2 ok
using K3 ok' Atrn2 eqSec-S-Cons trn24 trn34
unfolding statA' Van.completedFrom-def
using tr1' tr2' trn3 trn4 by force
next
assume K4: ?K4
then obtain sv1' sv2' statO' where 0:
statO' = sstatO' statO sv1 sv2
validTransV (sv1,sv1')
eqSec sv1 s1
validTransV (sv2,sv2')
eqSec sv2 s2
Van.eqAct sv1 sv2
and ok': statA' = Diff  $\longrightarrow$  statO' = Diff and  $\Delta$ :  $\Delta \infty \infty \infty s1' s2'$ 
statA' sv1' sv2' statO'
unfolding match12-12-def by auto
have r12': reachV sv1' reachV sv2' using r1 r2 0
by (metis Van.reach.Step fst-conv snd-conv)+
obtain trv1 trv2 where  $\psi$ :  $\psi s1' tr1' s2' tr2' statO' sv1' trv1 sv2' trv2$ 
using less(1)[of tr1' tr2', OF -  $\Delta$  - - - - - r34' r12', simplified]
using less.premis tr1' tr2' ok' A34' tr1'NE tr2'NE unfolding tr1 tr2
Opt.completedFrom-def by auto

```

show ?thesis **apply**(rule exI[of - sv1 # trv1]) **apply**(rule exI[of - sv2 # trv2])
using trn34
using ψ O33' O44' isAO34 tr1'NE tr2'NE **unfolding** ψ -def tr1 tr2
using Van.validFromS-Cons 0
using K4 eqSec-S-Cons
unfolding statA' Van.eqAct-def Van.completedFrom-def match12-12-def sstatO'-def
by simp (smt (z3) Simple-Transition-System.lastt-Cons Van.A.Cons-unfold Van.O.Cons-unfold list.inject status.exhaust status.simps(1) tr1' tr2' trn3 trn4 updStat.simps(4) updStat-diff)
qed
qed
qed
qed
qed

lemma unwindCond-final:
unwindCond $\Delta \implies reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$
 $(finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO\ s2)$
unfolding unwindCond-def
unfolding proact-def match-def match1-def match1-1-def
by auto

definition $\varphi\ \Delta\ w\ w1\ w2\ w1'\ w2'\ statA\ s1\ tr1\ s2\ tr2\ statAA\ statO\ sv1\ trv1\ sv2\ trv2\ statOO \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge$
 $(length\ trv1 > Suc\ 0 \vee w1' \leq w1) \wedge (length\ trv2 > Suc\ 0 \vee w2' \leq w2) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge$
 $Van.validFromS\ sv2\ trv2 \wedge$
 $Van.S\ trv1 = Opt.S\ tr1 \wedge Van.S\ trv2 = Opt.S\ tr2 \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $(statO = Eq \longrightarrow (statOO = Diff \longleftrightarrow Van.O\ trv1 \neq Van.O\ trv2)) \wedge$
 $(statA = Eq \longrightarrow (statAA = Diff \longleftrightarrow Opt.O\ tr1 \neq Opt.O\ tr2)) \wedge$

 $(statO = Diff \longrightarrow statOO = Diff) \wedge$
 $(statAA = Diff \longrightarrow statOO = Diff) \wedge$
 $\Delta\ w\ w1'\ w2'\ (lastt\ s1\ tr1)\ (lastt\ s2\ tr2)\ statAA\ (lastt\ sv1\ trv1)\ (lastt\ sv2\ trv2)$
statOO

lemma φ -final:
assumes unw: unwindCond Δ
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and vtr14: Opt.validFromS s1 tr1 Opt.validFromS s2 tr2

and $\varphi: \varphi \Delta w w1 w2 w1' w2' statA s1 tr1 s2 tr2 statAA statO sv1 trv1 sv2 trv2 statOO$
shows $(finalV (lastt sv1 trv1) \longleftrightarrow finalO (lastt s1 tr1)) \wedge (finalV (lastt sv2 trv2) \longleftrightarrow finalO (lastt s2 tr2))$
proof –
have $rsv12: Van.validFromS sv1 trv1 \longrightarrow reachV (lastt sv1 trv1)$
 $Van.validFromS sv2 trv2 \longrightarrow reachV (lastt sv2 trv2)$ **using** r
by $(simp\ add: Van.reach-validFromS-reach\ lastt-def)+$
have $rs14: Opt.validFromS s1 tr1 \longrightarrow reachO (lastt s1 tr1)$
 $Opt.validFromS s2 tr2 \longrightarrow reachO (lastt s2 tr2)$ **using** r
by $(simp\ add: Opt.reach-validFromS-reach\ lastt-def)+$
show $?thesis$ **using** $\varphi[unfolding\ \varphi-def]$ $rsv12\ rs14$ **using** $unw[unfolding\ unwind-Cond-def, rule-format,$
 $of\ lastt\ s1\ tr1\ lastt\ s2\ tr2\ lastt\ sv1\ trv1\ lastt\ sv2\ trv2\ w\ w1'\ w2'\ statAA\ statOO]$
using $vtr14(1)\ vtr14(2)$ **by** $auto$
qed

lemma φ -completedFrom: $unwindCond\ \Delta \implies$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $Opt.validFromS\ s1\ tr1 \implies completedFromO\ s1\ tr1 \implies$
 $Opt.validFromS\ s2\ tr2 \implies completedFromO\ s2\ tr2 \implies$
 $\varphi\ \Delta\ statA\ w\ w1\ w2\ w1'\ w2'\ s1\ tr1\ s2\ tr2\ statAA\ statO\ sv1\ trv1\ sv2\ trv2\ statOO$
 $\implies completedFromV\ sv1\ trv1 \wedge completedFromV\ sv2\ trv2$
using φ -final
by $(metis\ Van.completedFrom-def\ completedFromO-lastt\ lastt-def)$

lemma $unwindCond$ -ex- φ :
assumes $unwind: unwindCond\ \Delta$
and $\Delta: \Delta w w1 w2 s1 s2 statA sv1 sv2 statO$
and $r: reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$
and $stat: (statA = Diff \longrightarrow statO = Diff)$
and $v: Opt.validFromS\ s1\ tr1\ Opt.validFromS\ s2\ tr2$
and $i: isIntO (lastt\ s1\ tr1)\ isIntO (lastt\ s2\ tr2)$
and $nev: never\ isIntO (butlast\ tr1)\ never\ isIntO (butlast\ tr2)$
shows $\exists w' w1' w2' trv1 trv2 statAA statOO. \varphi\ \Delta\ w'\ w1\ w2\ w1'\ w2'\ statA\ s1\ tr1\ s2\ tr2\ statAA\ statO\ sv1\ trv1\ sv2\ trv2\ statOO$
using $assms(2-)$
proof $(induction\ length\ tr1 + length\ tr2\ w$
 $arbitrary: w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO\ tr1\ tr2\ rule: less2-induct')$
case $(less\ w\ tr1\ tr2\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO)$
note $ok = \langle statA = Diff \longrightarrow statO = Diff \rangle$
note $\Delta = \langle \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \rangle$
note $r34 = less(4,5)$ **note** $r12 = less(6,7)$
note $r = r34\ r12$
note $r3 = r34(1)$ **note** $r4 = r34(2)$ **note** $r1 = r12(1)$ **note** $r2 = r12(2)$
note $nev34 = less(13,14)$
note $nev3 = nev34(1)$ **note** $nev4 = nev34(2)$

have $i34: statA = Eq \longrightarrow isIntO\ s1 = isIntO\ s2$

and $f34$: $finalO\ s1 = finalO\ s2 \wedge finalV\ sv1 = finalO\ s1 \wedge finalV\ sv2 = finalO\ s2$
using Δ *unwind*[*unfolded unwindCond-def*] r **by** *auto*

note $is1 = \langle isIntO\ (lastt\ s1\ tr1) \rangle$
note $is2 = \langle isIntO\ (lastt\ s2\ tr2) \rangle$
note $utr1 = \langle Opt.validFromS\ s1\ tr1 \rangle$
note $utr2 = \langle Opt.validFromS\ s2\ tr2 \rangle$

have *proact-match*: $(\exists v < w. proact\ \Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO) \vee match\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
using Δ *unwind*[*unfolded unwindCond-def*] r **by** *auto*
show *?case using proact-match proof safe*
fix v **assume** $v: v < w$
assume *proact* $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
thus *?thesis unfolding proact-def proof safe*
assume $sv1: \neg isSecV\ sv1 \neg isIntV\ sv1$ **and** *move-1* $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
then obtain $sv1'$
where $0: validTransV\ (sv1, sv1')$
and $\Delta: \Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1'\ sv2\ statO$
unfolding *move-1-def* **by** *auto*
have $r1'$: *reachV* $sv1'$ **using** $r1\ 0$ **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain $w'\ w1'\ w2'\ trv1\ trv2\ statAA\ statOO$ **where** $\varphi: \varphi\ \Delta\ w'\ w1\ w2\ w1'\ w2'\ statA\ s1\ tr1\ s2\ tr2\ statAA\ statO\ sv1'\ trv1\ sv2\ trv2\ statOO$
using *less(2)*[*OF* v , *of* $tr1\ tr2\ w1\ w2\ s1\ s2\ statA\ sv1'\ sv2\ statO$, *simplified*, *OF* $\Delta\ r34\ r1'\ r2\ ok$]
using $is1\ is2\ nev3\ nev4\ utr1\ utr2$ **by** *blast*
show *?thesis apply(rule exI[of - w']) apply(rule exI[of - w1'])*
apply(*rule exI[of - w2']*) **apply**(*rule exI[of - sv1 # trv1]*) **apply**(*rule exI[of - trv2]*)
using $\varphi\ ok\ 0\ sv1$ **unfolding** φ -*def* **by** *auto*
next
assume $sv2: \neg isSecV\ sv2 \neg isIntV\ sv2$ **and** *move-2* $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
then obtain $sv2'$
where $0: validTransV\ (sv2, sv2')$
and $\Delta: \Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2'\ statO$
unfolding *move-2-def* **by** *auto*
have $r2'$: *reachV* $sv2'$ **using** $r2\ 0$ **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain $w'\ w1'\ w2'\ trv1\ trv2\ statAA\ statOO$ **where** $\varphi: \varphi\ \Delta\ w'\ w1\ w2\ w1'\ w2'\ statA\ s1\ tr1\ s2\ tr2\ statAA\ statO\ sv1\ trv1\ sv2'\ trv2\ statOO$
using *less(2)*[*OF* v , *of* $tr1\ tr2\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2'\ statO$, *simplified*, *OF* $\Delta\ r34\ r1\ r2'\ ok$]
using $is1\ is2\ nev3\ nev4\ utr1\ utr2$ **by** *blast*
show *?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule exI[of - w2'])*
apply(*rule exI[of - trv1]*) **apply**(*rule exI[of - sv2 # trv2]*)
using $\varphi\ ok\ 0\ sv2$ **unfolding** φ -*def* **by** *auto*

```

next
  assume sv12:  $\neg$  isSecV sv1  $\neg$  isSecV sv2 Van.eqAct sv1 sv2
  and move-12  $\Delta$  v w1 w2 s1 s2 statA sv1 sv2 statO
  then obtain sv1' sv2' statO'
  where 0: statO' = sstatO' statO sv1 sv2
  validTransV (sv1,sv1')  $\neg$  isSecV sv1
  validTransV (sv2,sv2')  $\neg$  isSecV sv2
  Van.eqAct sv1 sv2
  and  $\Delta$ :  $\Delta$  v w1 w2 s1 s2 statA sv1' sv2' statO'
  unfolding move-12-def by auto
  have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
fst-conv snd-conv)+
  have ok': statA = Diff  $\longrightarrow$  statO' = Diff
  using ok 0 unfolding sstatO'-def by (cases statO, auto)
  obtain w' w1' w2' trv1 trv2 statAA statOO where  $\varphi$ :  $\varphi$   $\Delta$  w' w1 w2 w1' w2'
statA s1 tr1 s2 tr2 statAA statO' sv1' trv1 sv2' trv2 statOO
  using less(2)[OF v, of tr1 tr2 w1 w2 s1 s2 statA sv1' sv2' statO', simplified,
OF  $\Delta$  r34 r12' ok']
  using is1 is2 nev3 nev4 vtr1 vtr2 by blast
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2'])
  apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
  apply(rule exI[of - statAA]) apply(rule exI[of - statOO])
  using  $\varphi$  ok' 0 sv12 nev unfolding  $\varphi$ -def sstatO'-def
  by simp (smt (verit, ccfv-SIG) Statewise-Attacker-Mod.eqAct-def
Van.A.Cons-unfold Van.O.Cons-unfold Van.Statewise-Attacker-Mod-axioms
Van.validFromS-Cons list.inject updStat.simps(1) updStat.simps(4))
qed
next
  assume m: match  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO
  define statA' where statA': statA' = sstatA' statA s1 s2
  show ?thesis
  proof(cases length tr1  $\leq$  Suc 0)
  case True
  hence tr1e: tr1 = []  $\vee$  tr1 = [s1]
  by (metis Opt.validFromS-singl-iff Suc-length-conv le-Suc-eq le-zero-eq length-0-conv
vtr1)
  hence Opt.A tr1 = [] by (simp add: True)
  hence Opt.A tr2 = [] using Opt.A.eq-Nil-iff nev4 by blast
  show ?thesis
  proof(cases length tr2  $\leq$  Suc 0)
  case True
  hence tr2e: tr2 = []  $\vee$  tr2 = [s2]
  by (metis Opt.validFromS-def Suc-length-conv le-Suc-eq le-zero-eq length-0-conv
list.sel(1) vtr2)
  show ?thesis apply(rule exI[of - w]) apply(rule exI[of - w1]) apply(rule
exI[of - w2])
  apply(rule exI[of - [sv1]], rule exI[of - [sv2]], rule exI[of - statA], rule exI[of
- statO])

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using tr1e tr2e
using f34 Δ apply (clarsimp simp: φ-def lastt-def)
apply(cases statA, simp-all)
apply (metis Opt.O.simps(4) Opt.S.simps(4) last-ConsL)
by (metis Opt.S.simps(4) last.simps ok)
next
case False
then obtain s24 tr2' where tr2: tr2 = s24 # tr2' and tr2'NE: tr2' ≠ []
by (cases tr2, auto)
have s24[simp]: s24 = s2 using ⟨Opt.validFromS s2 tr2⟩
by (simp add: Opt.validFromS-Cons-iff tr2)
obtain s2' where
trn4: validTransO (s2,s2') ∨ (s2 = s2' ∧ tr2' = []) and
tr2': Opt.validFromS s2' tr2' using ⟨Opt.validFromS s2 tr2⟩
unfolding tr2 s24 using Opt.validFromS-Cons-iff by auto
have r4': reachO s2'
using r4 trn4 by (metis Opt.reach.Step fst-conv snd-conv)+
have nev4': never isIntO (butlast tr2')
by (metis Opt.O.Nil-iff Opt.O.eq-Nil-iff nev4 tr2)
have isAO4: ¬ isIntO s2
using ⟨Opt.A tr2 = []⟩ tr2 tr2'NE by auto
have O44': Opt.O tr2 = Opt.O tr2' Opt.A tr2 = Opt.A tr2'
using isAO4 ⟨Opt.A tr2 = []⟩ tr2 by auto
have m: match2 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
match-def by auto
have (∃ w1' < w1. ∃ w2' < w2. ¬ isSecO s2 ∧ Δ ∞ w1' w2' s1 s2' statA sv1
sv2 statO) ∨
(∃ w1' < w1. eqSec sv2 s2 ∧ ¬ isIntV sv2 ∧ match2-1 Δ w1' ∞ s1 s2 s2'
statA sv1 sv2 statO) ∨
(¬ isSecV sv1 ∧ eqSec sv2 s2 ∧ Van.eqAct sv1 sv2 ∧ match2-12 Δ ∞ ∞
s1 s2 s2' statA sv1 sv2 statO)
using isAO4 trn4 ok tr2'NE
using m[unfolded match2-def, rule-format, of s2'] by auto
thus ?thesis
proof safe
fix w1'' w2'' assume w12': w1'' < w1 w2'' < w2
assume ¬ isSecO s2 and Δ: Δ ∞ w1'' w2'' s1 s2' statA sv1 sv2 statO
hence S4: Opt.S tr2' = Opt.S tr2 unfolding tr2 by auto
obtain w' w1' w2' trv1 trv2 statAA statOO where φ: φ Δ w' w1'' w2''
w1' w2' statA s1 tr1 s2' tr2' statAA statO sv1 trv1 sv2 trv2 statOO
using less(1)[of tr1 tr2', OF - Δ r3 r4' - - - - - nev3 nev4', unfolded
tr2, simplified]
using is1 is2 vtr1 vtr2 tr2' ok tr2'NE trn4 r1 r2 tr2 by auto
show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
using φ O44' S4 tr2 tr2'NE trn4 tr2' w12' unfolding φ-def by auto
next
fix w1'' assume w1': w1'' < w1
assume trn24: eqSec sv2 s2 and

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Atrn2:  $\neg$  isIntV sv2 and match2-1  $\Delta$  w1''  $\infty$  s1 s2 s2' statA sv1 sv2 statO
then obtain sv2' where trn2: validTransV (sv2,sv2') and
 $\Delta$ :  $\Delta$   $\infty$  w1''  $\infty$  s1 s2' statA sv1 sv2' statO
unfolding match2-1-def by auto
have r2': reachV sv2' using r2 trn2 by (metis Van.reach.Step fst-conv
snd-conv)
obtain w' w1' w2' trv1 trv2 statAA statOO where  $\varphi$ :  $\varphi$   $\Delta$  w' w1''  $\infty$  w1'
w2' statA s1 tr1 s2' tr2' statAA statO sv1 trv1 sv2' trv2 statOO
using less(1)[of tr1 tr2', OF -  $\Delta$  r3 r4' r1 r2' - - - - nev3 nev4', unfolded
tr2, simplified]
using is1 is2 tr2' tr2 vtr1 ok tr2'NE trn4 by auto
show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using  $\varphi$  tr2'NE
using Van.validFromS-Cons trn2
using isAO4 ok Atrn2 eqSec-S-Cons trn24 tr2' trn4 w1'
unfolding  $\varphi$ -def tr2 s24
by auto
next
assume sv1:  $\neg$  isSecV sv1 and trn24: eqSec sv2 s2 and
Atrn12: Van.eqAct sv1 sv2 and match2-12  $\Delta$   $\infty$   $\infty$  s1 s2 s2' statA sv1 sv2
statO
then obtain sv1' sv2' statO' where
statO': statO' = sstatO' statO sv1 sv2 and
trn1: validTransV (sv1,sv1') and
trn2: validTransV (sv2,sv2') and
 $\Delta$ :  $\Delta$   $\infty$   $\infty$   $\infty$  s1 s2' statA sv1' sv2' statO'
unfolding match2-12-def by auto
have r12': reachV sv1' reachV sv2'
using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
obtain w' w1' w2' trv1 trv2 statAA statOO where  $\varphi$ :  $\varphi$   $\Delta$  w'  $\infty$   $\infty$  w1'
w2' statA s1 tr1 s2' tr2' statAA statO' sv1' trv1 sv2' trv2 statOO
using less(1)[of tr1 tr2', OF -  $\Delta$  r3 r4' r12' - - - - nev3 nev4', simplified]
using is1 is2 vtr1 tr2 tr2' ok tr2'NE trn4 unfolding tr2 statO' sstatO'-def
by auto
show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
using  $\varphi$  O44' tr2'NE sv1
using Van.validFromS-Cons trn1 trn2
using isAO4 ok Atrn12 eqSec-S-Cons trn24 tr2' trn4
unfolding  $\varphi$ -def tr2 Van.completedFrom-def Van.eqAct-def statO' sstatO'-def

by simp (smt (verit, ccfv-threshold) Van.A.Cons-unfold i34 is1 last-ConsL
lastt-def status.exhaust tr1e updStat.simps(2))
qed
qed
next
case False
then obtain s13 tr1' where tr1: tr1 = s13 # tr1' and tr1'NE: tr1'  $\neq$  []

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    by (cases tr1, auto)
  have s13[simp]: s13 = s1 using ⟨Opt.validFromS s1 tr1⟩
    by (simp add: Opt.validFromS-Cons-iff tr1)
  obtain s1' where
    trn3: validTransO (s1,s1') and
    tr1': Opt.validFromS s1' tr1' using ⟨Opt.validFromS s1 tr1⟩
  unfolding tr1 s13 by (metis tr1'NE Simple-Transition-System.validFromS-Cons-iff)
  have r3': reachO s1' using r3 trn3 by (metis Opt.reach.Step fst-conv snd-conv)
  have f3: ¬ finalO s1 using Opt.final-def trn3 by blast
  hence f4: ¬ finalO s2 using f34 by blast
  have nev3': never isIntO (butlast tr1')
  using nev3 tr1 tr1'NE by auto
  have isAO3: ¬ isIntO s1 using less.premis(11) tr1 tr1'NE by auto
  have O33': Opt.O tr1 = Opt.O tr1' Opt.A tr1 = Opt.A tr1'
  using isAO3 unfolding tr1 by auto
  have m: match1 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
  match-def by auto
  have (∃ w1' < w1. ∃ w2' < w2. ¬ isSecO s1 ∧ Δ ∞ w1' w2' s1' s2 statA sv1
  sv2 statO) ∨
    (∃ w2' < w2. eqSec sv1 s1 ∧ ¬ isIntV sv1 ∧ match1-1 Δ ∞ w2' s1 s1' s2
  statA sv1 sv2 statO) ∨
    (eqSec sv1 s1 ∧ ¬ isSecV sv2 ∧ Van.eqAct sv1 sv2 ∧ match1-12 Δ ∞ ∞
  s1 s1' s2 statA sv1 sv2 statO)
  using m isAO3 trn3 ok unfolding match1-def by auto
  thus ?thesis
  proof safe
    fix w1'' w2'' assume w12': w1'' < w1 w2'' < w2
    assume ¬ isSecO s1 and Δ: Δ ∞ w1'' w2'' s1' s2 statA sv1 sv2 statO
    hence S3: Opt.S tr1' = Opt.S tr1 unfolding tr1 by auto
    obtain w' w1' w2' trv1 trv2 statAA statOO where φ: φ Δ w' w1'' w2'' w1'
  w2' statA s1' tr1' s2 tr2 statAA statO sv1 trv1 sv2 trv2 statOO
    using less(1)[of tr1' tr2, OF - Δ r3' r4 r12, unfolded O33', simplified]
    using is1 is2 tr1' ok f3 f4 tr1'NE trn3 O33'(1) nev3' nev4 vtr2 unfolding
  tr1 by auto
    show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
  exI[of - w2']) apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
    using φ O33' S3 tr1 tr1'NE tr1' trn3 w12' unfolding φ-def by auto
  next
    fix w2'' assume w2': w2'' < w2
    assume trn13: eqSec sv1 s1 and
    Atrn1: ¬ isIntV sv1 and match1-1 Δ ∞ w2'' s1 s1' s2 statA sv1 sv2 statO
    then obtain sv1' where
    trn1: validTransV (sv1,sv1') and
    Δ: Δ ∞ ∞ w2'' s1' s2 statA sv1' sv2 statO
    unfolding match1-1-def by auto
    have r1': reachV sv1' using r1 trn1 by (metis Van.reach.Step fst-conv
  snd-conv)
    obtain w' w1' w2' trv1 trv2 statAA statOO where φ: φ Δ w' ∞ w2'' w1'
  w2' statA s1' tr1' s2 tr2 statAA statO sv1' trv1 sv2 trv2 statOO

```

```

using less(1)[of tr1' tr2, OF - Δ r3' r4 r1' r2, unfolded O33', simplified]
using is1 is2 tr1 nev3' nev4 vtr1 vtr2 tr1' ok f3 f4 tr1'NE trn3 O33'(1)
unfolding tr1 by auto
show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
using φ O33' unfolding φ-def tr1 Van.completedFrom-def
using Van.validFromS-Cons trn1 tr1'NE tr1' trn3
using isAO3 ok Atrn1 eqSec-S-Cons trn13 w2'
by auto
next
assume sv2: ¬ isSecV sv2 and trn13: eqSec sv1 s1 and
Atrn12: Van.eqAct sv1 sv2 and match1-12 Δ ∞ ∞ s1 s1' s2 statA sv1 sv2
statO
then obtain sv1' sv2' statO' where
statO': statO' = sstatO' statO sv1 sv2 and
trn1: validTransV (sv1, sv1') and
trn2: validTransV (sv2, sv2') and
Δ: Δ ∞ ∞ ∞ s1' s2 statA sv1' sv2' statO'
unfolding match1-12-def by auto
have r12': reachV sv1' reachV sv2'
using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
obtain w' w1' w2' trv1 trv2 statAA statOO where φ: φ Δ w' ∞ ∞ w1'
w2' statA s1' tr1' s2 tr2 statAA statO' sv1' trv1 sv2' trv2 statOO
using less(1)[of tr1' tr2, OF - Δ r3' r4 r12', unfolded O33', simplified]
using less.premis tr1' ok f3 f4 tr1'NE trn3 O33'(1) unfolding tr1 statO'
sstatO'-def by auto

have trv1NE: trv1 ≠ [] and trv2NE: trv2 ≠ [] using φ unfolding φ-def by
auto
have [simp]: Van.O (sv1 # trv1) = Van.O (sv2 # trv2) ↔ (isIntV sv1
→ getObsV sv1 = getObsV sv2) ∧ Van.O trv1 = Van.O trv2
using Atrn12 trv1NE trv2NE unfolding Van.O.map-filter Van.eqAct-def by
simp
show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
using φ O33' tr1'NE sv2
using Van.validFromS-Cons trn1 trn2
using isAO3 ok Atrn12 eqSec-S-Cons trn13 f3 f34 s13 tr1' trn3
unfolding φ-def tr1 Van.completedFrom-def Van.eqAct-def statO' sstatO'-def
apply clarsimp
by (smt (verit, ccfv-SIG) Van.A.Cons-unfold updStat.simps(1) updStat.simps(2)
updStat.simps(4))
qed
qed
qed
qed

```

definition φa Δ w w1 w2 w1' w2' statA s1 tr1 s2 tr2 statAA statO sv1 trv1 sv2 trv2 statOO ≡

$trv1 \neq [] \wedge trv2 \neq [] \wedge$
 $(length\ trv1 > Suc\ 0 \vee w1' < w1) \wedge (length\ trv2 > Suc\ 0 \vee w2' < w2) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge$
 $Van.validFromS\ sv2\ trv2 \wedge$
 $Van.S\ trv1 = Opt.S\ tr1 \wedge Van.S\ trv2 = Opt.S\ tr2 \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $(statO = Eq \longrightarrow (statOO = Diff \longleftrightarrow Van.O\ trv1 \neq Van.O\ trv2)) \wedge$
 $(statA = Eq \longrightarrow (statAA = Diff \longleftrightarrow Opt.O\ tr1 \neq Opt.O\ tr2)) \wedge$

 $(statO = Diff \longrightarrow statOO = Diff) \wedge$
 $(statAA = Diff \longrightarrow statOO = Diff) \wedge$
 $\Delta\ w\ w1'\ w2'\ (lastt\ s1\ tr1)\ (lastt\ s2\ tr2)\ statAA\ (lastt\ sv1\ trv1)\ (lastt\ sv2\ trv2)$
 $statOO$

lemma *unwindCond-ex- φ a-getActO*:

assumes *unwind*: *unwindCond* Δ

and Δ : $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and *r34*: *reachO* *s1* *reachO* *s2* **and** *r12*: *reachV* *sv1* *reachV* *sv2*

and *stat*: $(statA = Diff \longrightarrow statO = Diff)$

and *v*: *validTransO* (*s1*, *s1'*) *validTransO* (*s2*, *s2'*)

and *i34*: *isIntO* *s1* *isIntO* *s2* *getActO* *s1* = *getActO* *s2*

shows $\exists\ w1'\ w2'\ trv1\ trv2\ statOO.$

$\varphi_a\ \Delta\ \infty\ w1\ w2\ w1'\ w2'\ statA\ s1\ [s1, s1']\ s2\ [s2, s2']\ (sstatA'\ statA\ s1\ s2)$

statO *sv1* *trv1* *sv2* *trv2* *statOO*

using Δ *r12* *stat*

proof(*induction* *w* *arbitrary*: *w1* *w2* *sv1* *sv2* *statO* *rule*: *less-induct*)

case (*less* *w* *w1* *w2* *sv1* *sv2* *statO*)

note $\Delta = \langle \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \rangle$

note *r12* = *less.prem*s(2,3)

note *r1* = *r12*(1) **note** *r2* = *r12*(2)

note *r* = *r34* *r12*

note *stat* = $\langle statA = Diff \longrightarrow statO = Diff \rangle$

have *f34*: *finalO* *s1* = *finalO* *s2* \wedge *finalV* *sv1* = *finalO* *s1* \wedge *finalV* *sv2* = *finalO* *s2*

using Δ *unwind*[*unfolded* *unwindCond-def*] *r* **by** *auto*

have *proact-match*: $(\exists\ v < w. \text{proact}\ \Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO) \vee \text{match}$
 $\Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

using Δ *unwind*[*unfolded* *unwindCond-def*] *r* **by** *auto*

show *?case* **using** *proact-match* **proof** *safe*

fix *v* **assume** *v*: *v* < *w*

assume *proact* $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

thus *?thesis* **unfolding** *proact-def* **proof** *safe*

assume *sv1*: $\neg\ isSecV\ sv1 \neg\ isIntV\ sv1$ **and** *move-1* $\Delta\ v\ w1\ w2\ s1\ s2\ statA$
sv1 *sv2* *statO*

then obtain *sv1'*

where *0*: *validTransV* (*sv1*, *sv1'*)

and Δ : $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1'\ sv2\ statO$


```

unfolding move-1-def by auto
have  $r1'$ : reachV sv1' using  $r1\ 0$  by (metis Van.reach.Step fst-conv snd-conv)
obtain  $w1'\ w2'\ trv1\ trv2\ statOO$  where
 $\varphi$ :  $\varphi a\ \Delta\ \infty\ w1\ w2\ w1'\ w2'\ statA\ s1\ [s1,\ s1']\ s2\ [s2,\ s2']\ (sstatA'\ statA\ s1\ s2)$ 
statO sv1' trv1 sv2 trv2 statOO
using less(1)[OF v Δ r1' r2 stat] by auto
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
using  $\varphi\ 0\ sv1$  unfolding  $\varphi a$ -def apply simp
by (metis Van.validFromS-Cons)
next
assume  $sv2$ :  $\neg\ isSecV\ sv2\ \neg\ isIntV\ sv2$  and  $move-2\ \Delta\ v\ w1\ w2\ s1\ s2\ statA$ 
sv1 sv2 statO
then obtain  $sv2'$ 
where  $0$ : validTransV (sv2,sv2')
and  $\Delta$ :  $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2'\ statO$ 
unfolding move-2-def by auto
have  $r2'$ : reachV sv2' using  $r2\ 0$  by (metis Van.reach.Step fst-conv snd-conv)
obtain  $w1'\ w2'\ trv1\ trv2\ statOO$  where
 $\varphi$ :  $\varphi a\ \Delta\ \infty\ w1\ w2\ w1'\ w2'\ statA\ s1\ [s1,\ s1']\ s2\ [s2,\ s2']\ (sstatA'\ statA\ s1\ s2)$ 
statO sv1 trv1 sv2' trv2 statOO
using less(1)[OF v Δ r1 r2' stat] by auto
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using  $\varphi\ 0\ sv2$  unfolding  $\varphi a$ -def apply simp by (metis Van.validFromS-Cons)
next
assume  $sv12$ :  $\neg\ isSecV\ sv1\ \neg\ isSecV\ sv2\ Van.eqAct\ sv1\ sv2$ 
and  $move-12\ \Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ 
then obtain  $sv1'\ sv2'\ statO'$ 
where  $0$ :  $statO' = sstatO'\ statO\ sv1\ sv2$ 
validTransV (sv1,sv1')  $\neg\ isSecV\ sv1$ 
validTransV (sv2,sv2')  $\neg\ isSecV\ sv2$ 
Van.eqAct sv1 sv2
and  $\Delta$ :  $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1'\ sv2'\ statO'$ 
unfolding move-12-def by auto
have  $r12'$ : reachV sv1' reachV sv2' using  $r1\ r2\ 0$  by (metis Van.reach.Step
fst-conv snd-conv)+
have  $stat'$ :  $statA = Diff\ \longrightarrow\ statO' = Diff$ 
using  $stat\ 0$  unfolding sstatO'-def by (cases statO, auto)
obtain  $w1'\ w2'\ trv1\ trv2\ statOO$  where
 $\varphi$ :  $\varphi a\ \Delta\ \infty\ w1\ w2\ w1'\ w2'\ statA\ s1\ [s1,\ s1']\ s2\ [s2,\ s2']\ (sstatA'\ statA\ s1\ s2)$ 
statO' sv1' trv1 sv2' trv2 statOO
using less(1)[OF v Δ r12' stat'] unfolding  $\varphi a$ -def apply simp by metis
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
using  $\varphi\ 0$  unfolding  $\varphi a$ -def sstatO'-def apply clarsimp apply(intro conjI)
subgoal by auto
subgoal by auto
subgoal by (metis Van.A.Cons-unfold Van.eqAct-def)

```

```

    subgoal apply(rule exI[of - statOO]) apply simp
    by (smt (verit, ccfv-threshold) Van.O.Cons-unfold Van.eqAct-def
        list.inject updStat.simps(1) updStat.simps(3)) .
qed
next
  assume m: match  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO
  define statA' where statA': statA' = sstatA' statA s1 s2
  have m: match12  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
  match-def by auto
  have ( $\exists w1' w2'. w1' < w1 \wedge w2' < w2 \wedge \neg isSecO s1 \wedge \neg isSecO s2 \wedge (statA
  = statA' \vee statO = Diff) \wedge$ 
     $\Delta \infty w1' w2' s1' s2' statA' sv1 sv2 statO$ )
     $\vee$ 
    ( $\exists w2' < w2. \neg isSecO s2 \wedge$ 
      eqSec sv1 s1  $\wedge \neg isIntV sv1 \wedge (statA = statA' \vee statO = Diff) \wedge$ 
      match12-1  $\Delta \infty w2' s1' s2' statA' sv1 sv2 statO$ )
     $\vee$ 
    ( $\exists w1' < w1. \neg isSecO s1 \wedge$ 
      eqSec sv2 s2  $\wedge \neg isIntV sv2 \wedge (statA = statA' \vee statO = Diff) \wedge$ 
      match12-2  $\Delta \infty w1' s1' s2' statA' sv1 sv2 statO$ )
     $\vee$ 
    (eqSec sv1 s1  $\wedge$  eqSec sv2 s2  $\wedge$  Van.eqAct sv1 sv2  $\wedge$ 
      match12-12  $\Delta \infty s1' s2' statA' sv1 sv2 statO$ )
  using m unfolding match12-def
  by (simp add: Opt.eqAct-def i34(1) i34(2) i34(3) statA' v(1) v(2))
  thus ?thesis
  apply(elim disjE exE)
  subgoal for w1' w2' apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - [sv1]]) apply(rule exI[of - [sv2]])
  apply(rule exI[of - statO])
  using stat unfolding  $\varphi a$ -def statA'
  by (auto simp add: i34(1) i34(2) sstatA'-def lastt-def)
  subgoal for w2' apply(rule exI[of -  $\infty$ ]) apply(rule exI[of - w2'])
  unfolding match12-1-def apply(elim conjE exE) subgoal for sv1'
  apply(rule exI[of - [sv1, sv1']]) apply(rule exI[of - [sv2]])
  apply(rule exI[of - statO])
  using stat unfolding  $\varphi a$ -def statA'
  by (auto simp add: i34(1) i34(2) sstatA'-def lastt-def) .
  subgoal for w1' apply(rule exI[of - w1']) apply(rule exI[of -  $\infty$ ])
  unfolding match12-2-def apply(elim conjE exE) subgoal for sv2'
  apply(rule exI[of - [sv1]]) apply(rule exI[of - [sv2, sv2']])
  apply(rule exI[of - statO])
  using stat unfolding  $\varphi a$ -def statA'
  by (auto simp add: i34(1) i34(2) sstatA'-def lastt-def) .
  subgoal unfolding match12-12-def apply(elim conjE exE) subgoal for sv1'
  sv2'
  apply(rule exI[of -  $\infty$ ]) apply(rule exI[of -  $\infty$ ])
  apply(rule exI[of - [sv1, sv1']]) apply(rule exI[of - [sv2, sv2']])
  apply(rule exI[of - sstatO' statO sv1 sv2])

```

using *stat unfolding* $\varphi a\text{-def}$ *statA'*
by (*auto simp add: i34 i34 sstatA'-def sstatO'-def lastt-def Van.eqAct-def*) . .
qed
qed

lemma *unwindCond-ex- $\varphi a'$ -aux:*
assumes *unwind: unwindCond* Δ
and $\Delta: \Delta$ *w w1 w2 s1 s2 statA sv1 sv2 statO*
and *r: reachO s1 reachO s2 reachV sv1 reachV sv2*
and *stat: (statA = Diff \longrightarrow statO = Diff)*
and *tr14NE: tr1 \neq [] tr2 \neq []*
and *v3': Opt.validFromS s1 (tr1 ## s1')* **and** *v4': Opt.validFromS s2 (tr2 ## s2')*
and *i: isIntO (lastt s1 tr1) isIntO (lastt s2 tr2)*
and *A34: getActO (lastt s1 tr1) = getActO (lastt s2 tr2)*
and *nev: never isIntO (butlast tr1) never isIntO (butlast tr2)*
shows $\exists w1' w2' trv1' trv2' statAA' statOO'$.
 $\varphi a \Delta \infty w1 w2 w1' w2' statA s1 (tr1 ## s1') s2 (tr2 ## s2') statAA' statO$
 $sv1 trv1' sv2 trv2' statOO'$

proof–

have *v3: Opt.validFromS s1 tr1* **and** *s13': validTransO (lastt s1 tr1, s1')*
apply (*metis v3' Opt.validFromS-def Opt.validS-append1 Nil-is-append-conv hd-append2*)
by (*metis Opt.validFromS-def Opt.validS-validTrans append-is-Nil-conv lastt-def*
list.distinct(1) list.sel(1) tr14NE(1) v3')
have *v4': Opt.validFromS s2 tr2* **and** *s24': validTransO (lastt s2 tr2, s2')*
apply (*metis v4' Opt.validFromS-def Opt.validS-append1 Nil-is-append-conv hd-append2*)
by (*metis Opt.validFromS-def Opt.validS-validTrans append-is-Nil-conv lastt-def*
list.sel(1) list.simps(3) tr14NE(2) v4')

obtain *ww ww1 ww2 trv1 trv2 statAA statOO* **where** $\varphi: \varphi \Delta ww w1 w2 ww1$
 $ww2 statA s1 tr1 s2 tr2 statAA statO sv1 trv1 sv2 trv2 statOO$
using *unwindCond-ex- φ [OF unwind Δ r stat v3 v4 i nev]* **by** *auto*

have *trv12NE: trv1 \neq [] trv2 \neq []* **using** φ **unfolding** $\varphi\text{-def}$ **by** *auto*

define *ss1 ss2 ssv1 ssv2* **where** *ss1: ss1 \equiv lastt s1 tr1* **and** *ss2: ss2 \equiv lastt s2*
tr2

and *ssv1: ssv1 \equiv lastt sv1 trv1* **and** *ssv2: ssv2 \equiv lastt sv2 trv2*

have *ss1l: ss1 = last tr1* **by** (*simp add: lastt-def ss1 tr14NE(1)*)
have *tr1l: tr1 = butlast tr1 @ [ss1]* **by** (*simp add: ss1l tr14NE(1)*)
have *ss2l: ss2 = last tr2* **by** (*simp add: lastt-def ss2 tr14NE(2)*)
have *tr2l: tr2 = butlast tr2 @ [ss2]* **by** (*simp add: ss2l tr14NE(2)*)
have *ssv1l: ssv1 = last trv1* **using** φ **unfolding** $\varphi\text{-def}$ **by** (*metis lastt-def ssv1*)
have *trv1l: trv1 = butlast trv1 @ [ssv1]* **by** (*simp add: ssv1l trv12NE(1)*)
have *ssv2l: ssv2 = last trv2* **using** φ **unfolding** $\varphi\text{-def}$ **by** (*metis lastt-def ssv2*)
have *trv2l: trv2 = butlast trv2 @ [ssv2]* **by** (*simp add: ssv2l trv12NE(2)*)

have *iss14[*simp*]: isIntO ss1 isIntO ss2* **using** *i unfolding* *ss1 ss2* **by** *auto*

```

have giss14[simp]: getActO ss1 = getActO ss2
  using A34 ss1 ss2 by fastforce

have [simp]: Opt.O (tr1 ## s1') = Opt.O tr1 ## getObsO ss1
by (metis Opt.O-def <isIntO ss1> holds-filtermap-RCons snoc-eq-iff-butlast tr1l)
have [simp]: Opt.O (tr2 ## s2') = Opt.O tr2 ## getObsO ss2
by (metis Opt.O-def <isIntO ss2> holds-filtermap-RCons snoc-eq-iff-butlast tr2l)

have [simp]: Opt.A (tr1 ## s1') = Opt.A tr1 ## getActO ss1
by (metis Opt.A-def <isIntO ss1> holds-filtermap-RCons snoc-eq-iff-butlast tr1l)
have [simp]: Opt.A (tr2 ## s2') = Opt.A tr2 ## getActO ss2
by (metis Opt.A-def <isIntO ss2> holds-filtermap-RCons snoc-eq-iff-butlast tr2l)
have [simp]: Opt.A (tr1 ## s1') = Opt.A (tr2 ## s2')  $\longleftrightarrow$  Opt.A tr1 = Opt.A tr2
by simp

have rss: reachO ss1 reachO ss2 reachV ssv1 reachV ssv2
using Opt.reach-validFromS-reach r ss1 tr14NE(1) v3 apply blast
using Opt.reach-validFromS-reach r(2) ss2l tr14NE(2) v4 apply blast
using Van.reach-validFromS-reach  $\varphi$ -def  $\varphi$  r(3) ssv1l
apply (smt (verit, del-insts))
using Van.reach-validFromS-reach  $\varphi$ -def  $\varphi$  r(4) ssv2l
apply (smt (verit, del-insts)) .

have stat: statAA = Diff  $\longrightarrow$  statOO = Diff
and  $\Delta$ :  $\Delta$  ww ww1 ww2 ss1 ss2 statAA ssv1 ssv2 statOO
using  $\varphi$  unfolding  $\varphi$ -def ss1[symmetric] ss2[symmetric] ssv1[symmetric] ssv2[symmetric]
by auto

note vs13 = s13'[unfolded ss1[symmetric]] note vs24 = s24'[unfolded ss2[symmetric]]
have  $\exists$  w1' w2' trv1' trv2' statA' statO'
 $\varphi$   $\Delta$   $\infty$  ww1 ww2 w1' w2' statAA ss1 [ss1,s1'] ss2 [ss2,s2'] (sstata' statAA ss1
ss2) statOO ssv1 trv1' ssv2 trv2' statO'
using unwindCond-ex- $\varphi$ a-getActO[OF unwind  $\Delta$  rss stat vs13 vs24 iss14 giss14]
by blast

then obtain w1' w2' trv1' trv2' statA' statO' where
 $\varphi$ 1:  $\varphi$   $\Delta$   $\infty$  ww1 ww2 w1' w2' statAA ss1 [ss1,s1'] ss2 [ss2,s2'] statA' statOO
ssv1 trv1' ssv2 trv2' statO' by auto

have trv12'NE: trv1'  $\neq$  [] trv2'  $\neq$  [] using  $\varphi$ 1 unfolding  $\varphi$ a-def by auto

have [simp]: Van.O (butlast trv1 @ trv1') = Van.O trv1 @ Van.O trv1'
using trv12'NE unfolding  $\varphi$ -def Van.O.map-filter Opt.O.map-filter apply(subst
butlast-append) by simp

have [simp]: Van.O (butlast trv2 @ trv2') = Van.O trv2 @ Van.O trv2'
using trv12'NE unfolding  $\varphi$ -def Van.O.map-filter Opt.O.map-filter apply(subst
butlast-append) by simp

```

have $\text{Van.A } trv1' = \text{Van.A } trv2'$ **using** $\varphi 1$ **unfolding** $\varphi a\text{-def}$ **by** *auto*
moreover have $\text{length } (\text{Van.O } trv1') = \text{length } (\text{Van.A } trv1') \wedge \text{length } (\text{Van.O } trv2') = \text{length } (\text{Van.A } trv2')$
unfolding $\text{Van.A.map-filter } \text{Van.O.map-filter}$ **by** *auto*
ultimately have $\text{length } (\text{Van.O } trv1') = \text{length } (\text{Van.O } trv2')$ **by** *auto*
hence [*simp*]: $\text{Van.O } trv1 @ \text{Van.O } trv1' = \text{Van.O } trv2 @ \text{Van.O } trv2' \longleftrightarrow \text{Van.O } trv1 = \text{Van.O } trv2 \wedge \text{Van.O } trv1' = \text{Van.O } trv2'$ **by** *auto*

have $\text{len: } trv1 \neq [] \wedge trv2 \neq [] \wedge trv1' \neq [] \wedge trv2' \neq [] \wedge$
 $(\text{Suc } 0 < \text{length } trv1 \vee ww1 \leq w1) \wedge$
 $(\text{Suc } 0 < \text{length } trv1' \vee w1' < ww1) \wedge$
 $(\text{Suc } 0 < \text{length } trv2 \vee ww2 \leq w2) \wedge$
 $(\text{Suc } 0 < \text{length } trv2' \vee w2' < ww2)$
using $\varphi \varphi 1$ **unfolding** $\varphi\text{-def}$ $\varphi a\text{-def}$ **by** *auto*

show *?thesis*
apply(*rule* $\text{exI}[of - w1']$) **apply**(*rule* $\text{exI}[of - w2']$)
apply(*rule* $\text{exI}[of - \text{butlast } trv1 @ trv1']$) **apply**(*rule* $\text{exI}[of - \text{butlast } trv2 @ trv2']$)
apply(*rule* $\text{exI}[of - \text{statA}']$) **apply**(*rule* $\text{exI}[of - \text{statO}']$)
unfolding $\varphi a\text{-def}$ **apply**(*intro* *conjI*)
subgoal using $\varphi \varphi 1$ **unfolding** $\varphi\text{-def}$ $\varphi a\text{-def}$ **by** *auto*
subgoal using $\varphi \varphi 1$ **unfolding** $\varphi\text{-def}$ $\varphi a\text{-def}$ **by** *auto*
subgoal using *len*
by *simp* (*metis* *Suc-lessI* *add-is-1* *diff-is-0-eq* *length-greater-0-conv* *linorder-not-less*

order-trans *trans-less-add2*)
subgoal using *len*
by *simp* (*metis* *Suc-leI* *le-add-diff-inverse2* *length-greater-0-conv* *nless-le* *order-le-less-trans* *trans-less-add2*)
subgoal using $\varphi \varphi 1$ **unfolding** $\varphi\text{-def}$ $\varphi a\text{-def}$ *ssv1*
using *Van.validFromS-append* **by** *auto*
subgoal using $\varphi \varphi 1$ **unfolding** $\varphi\text{-def}$ $\varphi a\text{-def}$ *ssv2*
using *Van.validFromS-append* **by** *auto*
subgoal using $\varphi \varphi 1$ **unfolding** $\varphi\text{-def}$ $\varphi a\text{-def}$ *Van.S.map-filter* *Opt.S.map-filter*

apply(*subst* *tr1l*) **apply**(*subst* *butlast-append*) **by** *simp*
subgoal using $\varphi \varphi 1$ **unfolding** $\varphi\text{-def}$ $\varphi a\text{-def}$ *Van.S.map-filter* *Opt.S.map-filter*

apply(*subst* *tr2l*) **apply**(*subst* *butlast-append*) **by** *simp*
subgoal using $\varphi \varphi 1$ **unfolding** $\varphi\text{-def}$ $\varphi a\text{-def}$ *Van.A.map-filter* *Opt.A.map-filter*

apply(*subst* *trv1l*) **apply**(*subst* *trv2l*)
apply(*subst* *butlast-append*) **apply** *simp* **apply**(*subst* *butlast-append*) **by** *simp*
subgoal using $\varphi \varphi 1$ **unfolding** $\varphi\text{-def}$ $\varphi a\text{-def}$ **apply** *simp*
apply(*cases* *Opt.O tr1 = Opt.O tr2, simp-all*) **apply** *clarify*
using *status.exhaust* **by** (*metis* (*full-types*))+
subgoal using $\varphi \varphi 1$ **unfolding** $\varphi\text{-def}$ $\varphi a\text{-def}$ **apply** *simp*
apply(*cases* *Opt.O tr1 = Opt.O tr2, simp-all*) **apply** *clarify*

apply (*smt* (*verit*, *del-insts*) *status.exhaust*)
by (*metis* *Opt.O.eq-Nil-iff* *nev(1)* *nev(2)*)
subgoal using φ $\varphi 1$ **unfolding** φ -def φa -def **by** *simp*
subgoal using φ $\varphi 1$ **unfolding** φ -def φa -def **by** *simp*
subgoal using $\varphi 1$ *trv12'NE* *tr14NE* **unfolding** φ -def φa -def *lastt-def* **by** *simp*

·
qed

lemma *unwindCond-ex- φa -aux2*:
assumes *unwind*: *unwindCond* Δ
and Δ : Δ *w* *w1* *w2* *s1* *s2* *statA* *sv1* *sv2* *statO*
and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*
and *stat*: (*statA* = *Diff* \longrightarrow *statO* = *Diff*)
and *v3'*: *Opt.validFromS* *s1* (*tr1* @ [*s1'*, *s1''*]) **and** *v4'*: *Opt.validFromS* *s2* (*tr2* @ [*s2'*, *s2''*])
and *i*: *isIntO* *s1'* *isIntO* *s2'*
and *A34*: *getActO* *s1'* = *getActO* *s2'*
and *nev*: *never isIntO* *tr1* *never isIntO* *tr2*
shows \exists *w1'* *w2'* *trv1* *trv2* *statAA* *statOO*.
 φa Δ ∞ *w1* *w2* *w1'* *w2'* *statA* *s1* (*tr1* @ [*s1'*, *s1''*]) *s2* (*tr2* @ [*s2'*, *s2''*]) *statAA*
statO *sv1* *trv1* *sv2* *trv2* *statOO*

proof–

have *0*: *lastt* *s1* (*tr1* ## *s1'*) = *s1'* *lastt* *s2* (*tr2* ## *s2'*) = *s2'*
unfolding *lastt-def* **by** *auto*
show *?thesis*
apply(*rule* *unwindCond-ex- $\varphi a'$ -aux[OF* *unwind* Δ *r* *stat*, *of* *tr1* ## *s1'* *tr2* ## *s2'*, *unfolded* *0*, *simplified*)
using *assms* **by** *auto*
qed

lemma *lastt-snoc[simp]*: *lastt* *s1* (*tr1* @ [*s1''*]) = *s1''*
unfolding *lastt-def* **by** *auto*

lemma *lastt-snoc2[simp]*: *lastt* *s1* (*tr1* @ [*s1'*, *s1''*]) = *s1''*
unfolding *lastt-def* **by** *auto*

lemma *append-snoc2*: *tr1* @ [*s1'*, *s1''*] = (*tr1* ## *s1'*) ## *s1''*
by *auto*

definition φ' Δ *w1* *w2* *w1'* *w2'* *statA* *s1* *tr1* *s1'* *s1''* *s2* *tr2* *s2'* *s2''* *statAA* *statO*
sv1 *trv1* *sv1''* *sv2* *trv2* *sv2''* *statOO* \equiv
(*trv1* \neq [] \vee *w1' < w1*) \wedge (*trv2* \neq [] \vee *w2' < w2*) \wedge
Van.validFromS *sv1* (*trv1* ## *sv1''*) \wedge *Van.validFromS* *sv2* (*trv2* ## *sv2''*) \wedge
Van.S (*trv1* ## *sv1''*) = *Opt.S* ((*tr1* ## *s1'*) ## *s1''*) \wedge *Van.S* (*trv2* ## *sv2''*) = *Opt.S* ((*tr2* ## *s2'*) ## *s2''*) \wedge
Van.A (*trv1* ## *sv1''*) = *Van.A* (*trv2* ## *sv2''*) \wedge
(*statO* = *Eq* \longrightarrow (*statOO* = *Diff*) = (*Van.O* (*trv1* ## *sv1''*) \neq *Van.O* (*trv2* ## *sv2''*))) \wedge

$(statA = Eq \longrightarrow (statAA = Diff) = (Opt.O ((tr1 \#\# s1') \#\# s1'') \neq Opt.O ((tr2 \#\# s2') \#\# s2''))) \wedge$
 $(statO = Diff \longrightarrow statOO = Diff) \wedge (statAA = Diff \longrightarrow statOO = Diff) \wedge$
 $\Delta \infty w1' w2' s1'' s2'' statAA sv1'' sv2'' statOO$

proposition *unwindCond-ex- φ'* :

assumes *unwind*: *unwindCond* Δ **and** Δ : $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$
and *r*: *reachO* *s1 reachO s2 reachV sv1 reachV sv2*
and *stat*: *statA = Diff \longrightarrow statO = Diff*
and *v3'*: *Opt.validFromS s1 ((tr1 \#\# s1') \#\# s1'')* **and** *v4'*: *Opt.validFromS s2 ((tr2 \#\# s2') \#\# s2'')*
and *i*: *isIntO s1' isIntO s2'*
and *A34*: *getActO s1' = getActO s2'*
and *nev*: *never isIntO tr1 never isIntO tr2*
shows $\exists w1' w2' trv1 sv1'' trv2 sv2'' statAA statOO.$
 $\varphi' \Delta w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$
using *unwindCond-ex- φ -aux2*[*unfolded φ -def, unfolded lastt-snoc lastt-snoc2 append-snoc2, OF assms*]
unfolding *φ -def* **apply**(*elim exE*) **subgoal for** *w1' w2' trv1 trv2 statAA statOO*
apply(*cases trv1 rule: rev-cases*)
subgoal by *auto*
apply(*cases trv2 rule: rev-cases*)
subgoal by *auto*
subgoal unfolding *φ' -def* **apply** *simp* **by** *blast . .*

definition $\chi^3 \Delta w (w1::enat) w2 w1' w2' s1 tr1 s2 statAA sv1 trv1 sv2 trv2 statOO \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge (length\ trv2 > Suc\ 0 \vee w2' \leq w2) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge Van.validFromS\ sv2\ trv2 \wedge$
 $never\ isSecV\ (butlast\ trv1) \wedge$
 $isSecV\ (lastt\ sv1\ trv1) \wedge getSecV\ (lastt\ sv1\ trv1) = getSecO\ (lastt\ s1\ tr1) \wedge$
 $never\ isSecV\ (butlast\ trv2) \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $\Delta w w1' w2' (lastt\ s1\ tr1) s2 statAA (lastt\ sv1\ trv1) (lastt\ sv2\ trv2) statOO$

lemma *χ^3 -final*:

assumes *unw*: *unwindCond* Δ
and *r*: *reachO s1 reachO s2 reachV sv1 reachV sv2*
and *vtr1*: *Opt.validFromS s1 tr1*
and χ^3 : $\chi^3 \Delta w w1 w2 w1' w2' s1 tr1 s2 statAA sv1 trv1 sv2 trv2 statOO$
shows (*finalV (lastt sv1 trv1)*) \longleftrightarrow (*finalO (lastt s1 tr1)*) \wedge (*finalV (lastt sv2 trv2)*)
 \longleftrightarrow (*finalO s2*)
proof –
have *rsv12*: *Van.validFromS sv1 trv1 \longrightarrow reachV (lastt sv1 trv1)*
 $Van.validFromS sv2 trv2 \longrightarrow reachV (lastt sv2 trv2)$ **using** *r*

by (*simp add: Van.reach-validFromS-reach lastt-def*)
have $rs1: \text{Opt.validFromS } s1 \ tr1 \longrightarrow \text{reachO } (\text{lastt } s1 \ tr1)$
using r
by (*simp add: Opt.reach-validFromS-reach lastt-def*)
show *?thesis using* $\chi^3[\text{unfolded } \chi^3\text{-def}] \ rsv12 \ rs1$ **using** *unw[unfolded unwind-Cond-def, rule-format,*
of lastt s1 tr1 s2 lastt sv1 trv1 lastt sv2 trv2 w w1' w2' statAA statOO]
using $\text{vtr1 } \langle \text{reachO } s2 \rangle$ **by** *auto*
qed

lemma $\chi^3\text{-completedFrom: unwindCond } \Delta \Longrightarrow$
 $\text{reachO } s1 \Longrightarrow \text{reachO } s2 \Longrightarrow \text{reachV } sv1 \Longrightarrow \text{reachV } sv2 \Longrightarrow$
 $\text{Opt.validFromS } s1 \ tr1 \Longrightarrow \text{completedFromO } s1 \ tr1 \Longrightarrow$
 $\chi^3 \ \Delta \ w \ w1 \ w2 \ w1' \ w2' \ s1 \ tr1 \ s2 \ \text{statAA} \ sv1 \ trv1 \ sv2 \ trv2 \ \text{statOO}$
 $\Longrightarrow \text{completedFromV } sv1 \ trv1 \ \wedge \ \text{completedFromV } sv2 \ trv2$
by (*metis Van.final-not-isSec* $\chi^3\text{-def}$ $\chi^3\text{-final}$ $\text{completedFromO-lastt}$)

lemma *unwindCond-ex- χ^3 :*
assumes *unwind: unwindCond* Δ
and $\Delta: \Delta \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$
and $r: \text{reachO } s1 \ \text{reachO } s2 \ \text{reachV } sv1 \ \text{reachV } sv2$
and $\text{vtr1}: \text{Opt.validFromS } s1 \ tr1$
and $\text{nis1}: \neg \text{isIntO } s1$ **and** $\text{nis2}: \neg \text{isIntO } s2$
and $\text{inter3}: \text{never isIntO } tr1$
and $\text{sec}: \text{never isSecO } (\text{butlast } tr1) \ \text{isSecO } (\text{lastt } s1 \ tr1)$
shows $\exists w' \ w1' \ w2' \ trv1 \ trv2 \ \text{statOO}. \ \chi^3 \ \Delta \ w' \ w1 \ w2 \ w1' \ w2' \ s1 \ tr1 \ s2 \ \text{statA} \ sv1$
 $\text{trv1} \ sv2 \ trv2 \ \text{statOO}$
using *assms(2-)*
proof(*induction length tr1 w*
arbitrary: w1 w2 s1 s2 statA sv1 sv2 statO tr1 rule: less2-induct')
case (*less w tr1 w1 w2 s1 s2 statA sv1 sv2 statO*)
note $\text{vtr1} = \text{less}(8)$

note $\Delta = \langle \Delta \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \rangle$
note $\text{nis1} = \text{less}(9)$ **note** $\text{nis2} = \text{less}(10)$
note $\text{inter3} = \text{less}(11)$
note $\text{sec3} = \text{less}(12,13)$
note $r34 = \text{less.prem}(2,3)$ **note** $r12 = \text{less.prem}(4,5)$
note $r = r34 \ r12$
note $r3 = r34(1)$ **note** $r4 = r34(2)$ **note** $r1 = r12(1)$ **note** $r2 = r12(2)$

have $i34: \text{statA} = \text{Eq} \longrightarrow \text{isIntO } s1 = \text{isIntO } s2$
and $f34: \text{finalO } s1 = \text{finalO } s2 \ \wedge \ \text{finalV } sv1 = \text{finalO } s1 \ \wedge \ \text{finalV } sv2 = \text{finalO}$
 $s2$
using Δ *unwind[unfolded unwindCond-def]* r **by** *auto*

have *proact-match: ($\exists v < w. \text{proact } \Delta \ v \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$) \vee match*
 $\Delta \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$
using Δ *unwind[unfolded unwindCond-def]* r **by** *auto*


```

show ?case using proact-match proof safe
  fix v assume v: v < w
  assume proact  $\Delta$  v w1 w2 s1 s2 statA sv1 sv2 statO
  thus ?thesis unfolding proact-def proof safe
  assume sv1:  $\neg$  isSecV sv1  $\neg$  isIntV sv1 and move-1  $\Delta$  v w1 w2 s1 s2 statA
  sv1 sv2 statO
  then obtain sv1'
  where 0: validTransV (sv1,sv1')
  and  $\Delta$ :  $\Delta$  v w1 w2 s1 s2 statA sv1' sv2 statO
  unfolding move-1-def by auto
  have r1': reachV sv1' using r1 0 by (metis Van.reach.Step fst-conv snd-conv)
  obtain w' w1' w2' trv1 trv2 statOO where  $\chi^3$ :  $\chi^3$   $\Delta$  w' w1 w2 w1' w2' s1
  tr1 s2 statA sv1' trv1 sv2 trv2 statOO
  using less(2)[OF v, of tr1 w1 w2 s1 s2 statA sv1' sv2 statO,
    simplified, OF  $\Delta$  r34 r1' r2 vtr1 nis1 nis2 inter3 sec3] by auto
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
  exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
  using  $\chi^3$  0 sv1 unfolding  $\chi^3$ -def by auto
  next
  assume sv2:  $\neg$  isSecV sv2  $\neg$  isIntV sv2 and move-2  $\Delta$  v w1 w2 s1 s2 statA
  sv1 sv2 statO
  then obtain sv2'
  where 0: validTransV (sv2,sv2')
  and  $\Delta$ :  $\Delta$  v w1 w2 s1 s2 statA sv1 sv2' statO
  unfolding move-2-def by auto
  have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
  obtain w' w1' w2' trv1 trv2 statOO where  $\chi^3$ :  $\chi^3$   $\Delta$  w' w1 w2 w1' w2' s1
  tr1 s2 statA sv1 trv1 sv2' trv2 statOO
  using less(2)[OF v, of tr1 w1 w2 s1 s2 statA sv1 sv2' statO,
    simplified, OF  $\Delta$  r34 r1 r2' vtr1 nis1 nis2 inter3 sec3] by auto
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
  exI[of - w2']) apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
  using  $\chi^3$  0 sv2 unfolding  $\chi^3$ -def by auto
  next
  assume sv12:  $\neg$  isSecV sv1  $\neg$  isSecV sv2 Van.eqAct sv1 sv2
  and move-12  $\Delta$  v w1 w2 s1 s2 statA sv1 sv2 statO
  then obtain sv1' sv2' statO'
  where 0: statO' = sstatO' statO sv1 sv2
  validTransV (sv1,sv1')  $\neg$  isSecV sv1
  validTransV (sv2,sv2')  $\neg$  isSecV sv2
  Van.eqAct sv1 sv2
  and  $\Delta$ :  $\Delta$  v w1 w2 s1 s2 statA sv1' sv2' statO'
  unfolding move-12-def by auto
  have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
  fst-conv snd-conv)+

  obtain w' w1' w2' trv1 trv2 statOO where  $\chi^3$ :  $\chi^3$   $\Delta$  w' w1 w2 w1' w2' s1
  tr1 s2 statA sv1' trv1 sv2' trv2 statOO
  using less(2)[OF v, of tr1 w1 w2 s1 s2 statA sv1' sv2' statO',

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      simplified, OF  $\Delta$  r34 r12' vtr1 nis1 nis2 inter3 sec3] by auto
    show ?thesis apply(rule exI[of - w]) apply(rule exI[of - w1]) apply(rule
exI[of - w2]) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
    apply(rule exI[of - statOO])
    using  $\chi^3$  0 sv12 unfolding  $\chi^3$ -def sstatO'-def
    by (auto simp: Van.eqAct-def)
  qed
next
assume m: match  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO
define statA' where statA': statA' = sstatA' statA s1 s2
show ?thesis
proof(cases length tr1  $\leq$  Suc 0)
  case True
    hence tr1e: tr1 = []  $\vee$  tr1 = [s1]
  by (metis Opt.validFromS-singl-iff Suc-length-conv le-Suc-eq le-zero-eq length-0-conv
vtr1)
    hence Opt.A tr1 = [] by (simp add: True)
    have is1: isSecO s1
      by (metis last.simps lastt-def sec3(2) tr1e)
    hence  $\neg$  finalO s1 using Opt.final-not-isSec by blast
    then obtain s1' where s13': validTransO (s1, s1') unfolding Opt.final-def
  by auto
    hence isv1: isSecV sv1  $\wedge$  getSecV sv1 = getSecO s1 using m is1 nis1
    unfolding match-def match1-def eqSec-def by auto
    show ?thesis using tr1e isv1 apply-
      apply(rule exI[of - w]) apply(rule exI[of - w1]) apply(rule exI[of - w2])
      apply(rule exI[of - [sv1]], rule exI[of - [sv2]], rule exI[of - statO])
      using tr1e
      using f34  $\Delta$  by (clarsimp simp:  $\chi^3$ -def lastt-def)
  next
  case False
    then obtain s13 tr1' where tr1: tr1 = s13 # tr1' and tr1'NE: tr1'  $\neq$  []
      by (cases tr1, auto)
    have s13[simp]: s13 = s1 using  $\langle$ Opt.validFromS s1 tr1 $\rangle$ 
      by (simp add: Opt.validFromS-Cons-iff tr1)
    obtain s1' where
      trn3: validTransO (s1, s1') and
      tr1': Opt.validFromS s1' tr1' using  $\langle$ Opt.validFromS s1 tr1 $\rangle$ 
    unfolding tr1 s13 by (metis tr1'NE Simple-Transition-System.validFromS-Cons-iff)
    have r3': reachO s1' using r3 trn3 by (metis Opt.reach.Step fst-conv snd-conv)
    have f3:  $\neg$  finalO s1 using Opt.final-def trn3 by blast
    hence f4:  $\neg$  finalO s2 using f34 by blast
    have nev3': never isIntO tr1'
      using inter3 tr1 tr1'NE by auto
    have isAO3:  $\neg$  isIntO s1 by (simp add: nis1)
    have O33': Opt.O tr1 = Opt.O tr1' Opt.A tr1 = Opt.A tr1'
      using isAO3 unfolding tr1 by auto
    have m: match1  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
match-def by auto

```

```

have ( $\exists w1' < w1. \exists w2' < w2. \neg isSecO\ s1 \wedge \Delta \infty w1' w2' s1' s2\ statA\ sv1\ sv2\ statO$ )  $\vee$ 
( $\exists w2' < w2. eqSec\ sv1\ s1 \wedge \neg isIntV\ sv1 \wedge match1-1\ \Delta \infty w2' s1\ s1' s2\ statA\ sv1\ sv2\ statO$ )  $\vee$ 
( $eqSec\ sv1\ s1 \wedge \neg isSecV\ sv2 \wedge Van.eqAct\ sv1\ sv2 \wedge match1-12\ \Delta \infty \infty\ s1\ s1' s2\ statA\ sv1\ sv2\ statO$ )
using  $m\ isAO3\ trn3$  unfolding  $match1-def$  by  $auto$ 
thus  $?thesis$ 
proof  $safe$ 
fix  $w1'' w2''$  assume  $w12': w1'' < w1\ w2'' < w2$ 
assume  $\neg isSecO\ s1$  and  $\Delta: \Delta \infty w1'' w2'' s1' s2\ statA\ sv1\ sv2\ statO$ 
hence  $S3: Opt.S\ tr1' = Opt.S\ tr1$  unfolding  $tr1$  by  $auto$ 
obtain  $w' w1' w2' trv1\ trv2\ statOO$  where  $\chi3: \chi3\ \Delta\ w' w1'' w2'' w1' w2'$ 
 $s1' tr1' s2\ statA\ sv1\ trv1\ sv2\ trv2\ statOO$ 
using  $less(1)[of\ tr1', OF - \Delta\ r3' r4' r12 -]$  unfolding  $tr1$ 
by  $simp\ (metis\ Opt.S.eq-Nil-iff(2)\ S3\ Opt.validFromS-def\ \langle \neg isSecO\ s1 \rangle\ last.simps)$ 
 $lastt-def\ list-all-hd\ nev3'\ nis2\ s13\ sec3(1)\ sec3(2)\ tr1\ tr1'$ 
show  $?thesis$  apply( $rule\ exI[of - w']$ ) apply( $rule\ exI[of - w1']$ ) apply( $rule\ exI[of - w2']$ )
apply( $rule\ exI[of - trv1]$ ) apply( $rule\ exI[of - trv2]$ )
using  $\chi3\ O33'$  unfolding  $\chi3-def\ tr1\ Van.completedFrom-def$ 
using  $Van.validFromS-Cons\ tr1'NE\ tr1' trn3\ isAO3\ w12'$  by  $auto$ 
next
fix  $w2''$  assume  $w2': w2'' < w2$ 
assume  $trn13: eqSec\ sv1\ s1$  and
 $Atrn1: \neg isIntV\ sv1$  and  $match1-1\ \Delta \infty w2'' s1\ s1' s2\ statA\ sv1\ sv2\ statO$ 
then obtain  $sv1'$  where
 $trn1: validTransV\ (sv1, sv1')$  and
 $\Delta: \Delta \infty \infty w2'' s1' s2\ statA\ sv1' sv2\ statO$ 
unfolding  $match1-1-def$  by  $auto$ 
have  $r1': reachV\ sv1'$  using  $r1\ trn1$  by ( $metis\ Van.reach.Step\ fst-conv\ snd-conv$ )
obtain  $w' w1' w2' trv1\ trv2\ statOO$  where  $\chi3: \chi3\ \Delta\ w' \infty w2'' w1' w2'$ 
 $s1' tr1' s2\ statA\ sv1' trv1\ sv2\ trv2\ statOO$ 

using  $less(1)[of\ tr1', OF - \Delta\ r3' r4' r1' r2, unfolded\ O33', simplified]$ 
using  $less.prem\ tr1' f3\ f4\ tr1'NE\ trn3\ O33'(1)$ 
unfolding  $tr1$ 
by  $simp\ (metis\ Opt.validFromS-def\ list-all-hd)$ 
show  $?thesis$  apply( $rule\ exI[of - w']$ ) apply( $rule\ exI[of - w1']$ ) apply( $rule\ exI[of - w2']$ )
apply( $rule\ exI[of - sv1\ \# trv1]$ ) apply( $rule\ exI[of - trv2]$ )
using  $\chi3\ O33'$  unfolding  $\chi3-def\ tr1\ Van.completedFrom-def$ 
using  $Van.validFromS-Cons\ trn1\ tr1'NE\ tr1' trn3$ 
using  $isAO3\ Atrn1\ eqSec-S-Cons\ trn13\ w2'$ 
by  $simp\ (metis\ Opt.S.Nil-iff\ Opt.S.eq-Nil-iff(1)\ eqSec-def\ nless-le\ order-le-less-trans\ s13\ sec3(1)\ tr1)$ 
next
assume  $sv2: \neg isSecV\ sv2$  and  $trn13: eqSec\ sv1\ s1$  and
 $Atrn12: Van.eqAct\ sv1\ sv2$  and  $match1-12\ \Delta \infty \infty\ s1\ s1' s2\ statA\ sv1\ sv2$ 

```

```

statO
  then obtain sv1' sv2' statO' where
    statO': statO' = sstatO' statO sv1 sv2 and
    trn1: validTransV (sv1,sv1') and
    trn2: validTransV (sv2,sv2') and
    Δ: Δ ∞ ∞ ∞ s1' s2 statA sv1' sv2' statO'
  unfolding match1-12-def by auto
  have r12': reachV sv1' reachV sv2'
  using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
  obtain w' w1' w2' trv1 trv2 statOO where χ3: χ3 Δ w' ∞ ∞ w1' w2' s1'
tr1' s2 statA sv1' trv1 sv2' trv2 statOO
  using less(1)[of tr1', OF - Δ r3' r4' r12', unfolded O33', simplified]
  using less.premis tr1' f3 f4 tr1'NE trn3 O33'(1) unfolding tr1 statO'
sstatO'-def
  by simp (metis Simple-Transition-System.validFromS-def list-all-hd)+
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
  using χ3 O33' tr1'NE sv2
  using Van.validFromS-Cons trn1 trn2
  using isAO3 Atrn12 eqSec-S-Cons trn13 f3 f34 s13 tr1' trn3
  unfolding χ3-def tr1 Van.completedFrom-def Van.eqAct-def
  using Van.A.Cons-unfold eqSec-def sec3(1) tr1 by auto
qed
qed
qed
qed

```

definition $\chi3a$ where $\chi3a \Delta w (w1::enat) w2 w1' w2' s1 s1' s2 statAA sv1 trv1 sv2 trv2 statOO \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge (length\ trv2 > Suc\ 0 \vee w2' < w2) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge Van.validFromS\ sv2\ trv2 \wedge$
 $Van.S\ trv1 = [getSecO\ s1] \wedge$
 $never\ isSecV\ (butlast\ trv2) \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $\Delta\ w\ w1'\ w2'\ s1'\ s2\ statAA\ (lastt\ sv1\ trv1)\ (lastt\ sv2\ trv2)\ statOO$

lemma *unwindCond-ex-χ3a-getSec*:
assumes *unwind*: *unwindCond* Δ
and Δ: Δ w w1 w2 s1 s2 statA sv1 sv2 statO
and r34: reachO s1 reachO s2 **and** r12: reachV sv1 reachV sv2
and v: validTransO (s1, s1')
and ii3: ¬ isIntO s1
and is1: isSecO s1 **and** isv13: isSecV sv1 getSecO s1 = getSecV sv1
shows ∃ w1' w2' trv1 trv2 statOO.
 $\chi3a \Delta \infty w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv2 trv2 statOO$
using Δ r12 isv13
proof(*induction w arbitrary: w1 w2 sv1 sv2 statO rule: less-induct*)
case (*less w w1 w2 sv1 sv2 statO*)
note Δ = ⟨Δ w w1 w2 s1 s2 statA sv1 sv2 statO⟩

```

note r12 = less.premis(2,3)
note r1 = r12(1) note r2 = r12(2)
note r = r34 r12
note isv13 = ⟨isSecV sv1⟩ ⟨getSecO s1 = getSecV sv1⟩

have f34: finalO s1 = finalO s2 ∧ finalV sv1 = finalO s1 ∧ finalV sv2 = finalO
s2
  using Δ unwind[unfolded unwindCond-def] r by auto

have proact-match: (∃ v<w. proact Δ v w1 w2 s1 s2 statA sv1 sv2 statO) ∨ match
Δ w1 w2 s1 s2 statA sv1 sv2 statO
  using Δ unwind[unfolded unwindCond-def] r by auto
  show ?case using proact-match proof safe
    fix v assume v: v < w
    assume proact Δ v w1 w2 s1 s2 statA sv1 sv2 statO
    thus ?thesis unfolding proact-def proof safe
      assume sv1: ¬ isSecV sv1 ¬ isIntV sv1 and move-1 Δ v w1 w2 s1 s2 statA
sv1 sv2 statO
      hence False using isv13 by blast
      thus ?thesis by auto
    next
      assume sv2: ¬ isSecV sv2 ¬ isIntV sv2 and move-2 Δ v w1 w2 s1 s2 statA
sv1 sv2 statO
      then obtain sv2'
      where 0: validTransV (sv2,sv2')
      and Δ: Δ v w1 w2 s1 s2 statA sv1 sv2' statO
      unfolding move-2-def by auto
      have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
      obtain w1' w2' trv1 trv2 statOO where
      χ3a: χ3a Δ ∞ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv2' trv2 statOO
      using less(1)[OF v Δ r1 r2' isv13] by auto
      show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
      using χ3a 0 sv2 unfolding χ3a-def by auto
    next
      assume sv12: ¬ isSecV sv1 ¬ isSecV sv2 Van.eqAct sv1 sv2
      and move-12 Δ v w1 w2 s1 s2 statA sv1 sv2 statO
      hence False using isv13 by blast
      thus ?thesis by auto
    qed
  next
    assume m: match Δ w1 w2 s1 s2 statA sv1 sv2 statO
    have m: match1 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
match-def by auto
    have (∃ w1' w2'. w1'<w1 ∧ w2'<w2 ∧ ¬ isSecO s1 ∧ Δ ∞ w1' w2' s1' s2
statA sv1 sv2 statO) ∨
      (∃ w2'<w2. eqSec sv1 s1 ∧ ¬ isIntV sv1 ∧ match1-1 Δ ∞ w2' s1 s1' s2
statA sv1 sv2 statO) ∨
      (eqSec sv1 s1 ∧ ¬ isSecV sv2 ∧ Van.eqAct sv1 sv2 ∧ match1-12 Δ ∞ ∞

```

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s1 s1' s2 statA sv1 sv2 statO)
  using m v ii3 unfolding match1-def by auto

thus ?thesis
apply (elim disjE exE)
  subgoal for w1' w2' using is1 by auto
  subgoal for w2' apply (rule exI[of - ∞]) apply (rule exI[of - w2'])
  unfolding match1-1-def apply (elim conjE exE) subgoal for sv1'
  apply (rule exI[of - [sv1, sv1']]) apply (rule exI[of - [sv2]])
  apply (rule exI[of - statO])
  using is1 isv13 unfolding χ3a-def
  by (auto simp : sstatA'-def lastt-def) .
  subgoal apply (rule exI[of - ∞]) apply (rule exI[of - ∞])
  unfolding match1-12-def apply (elim conjE exE) subgoal for sv1' sv2'
  apply (rule exI[of - [sv1, sv1']]) apply (rule exI[of - [sv2, sv2']])
  apply (rule exI[of - sstatO' statO sv1 sv2])
  using is1 isv13 unfolding χ3a-def
  by (auto simp : sstatA'-def sstatO'-def lastt-def Van.eqAct-def) .
qed
qed

```

definition $\chi3b \Delta w (w1::enat) w2 w1' w2' s1 tr1 s2 statAA sv1 trv1 sv2 trv2$
 $statOO \equiv$
 $trv1 \neq [] \wedge$
 $trv2 \neq [] \wedge (length\ trv2 > Suc\ 0 \vee w2' < w2) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge$
 $Van.validFromS\ sv2\ trv2 \wedge$
 $Van.S\ trv1 = Opt.S\ tr1 \wedge$
 $never\ isSecV\ (butlast\ trv2) \wedge Van.A\ trv1 = Van.A\ trv2 \wedge$
 $\Delta\ w\ w1'\ w2'\ (lastt\ s1\ tr1)\ s2\ statAA\ (lastt\ sv1\ trv1)\ (lastt\ sv2\ trv2)\ statOO$

lemma *unwindCond-ex-χ3b-aux*:
assumes *unwind*: *unwindCond* Δ
and Δ : $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ **and**
 r : *reachO* $s1$ *reachO* $s2$ *reachV* $sv1$ *reachV* $sv2$
and *tr1NE*: $tr1 \neq []$
and $v3'$: *Opt.validFromS* $s1$ ($tr1 \## s1'$)
and $nis1$: $\neg isIntO\ s1$ **and** $nis2$: $\neg isIntO\ s2$
and $ninter3'$: *never isIntO* ($tr1 \## s1'$)
and sec : *never isSecO* (*butlast* $tr1$) *isSecO* (*lastt* $s1\ tr1$)
shows $\exists w1'\ w2'\ trv1\ trv2\ statOO. \chi3b\ \Delta\ \infty\ w1\ w2\ w1'\ w2'\ s1\ (tr1 \## s1')\ s2$
 $statA\ sv1\ trv1\ sv2\ trv2\ statOO$
proof –
have $v3$: *Opt.validFromS* $s1\ tr1$ **and** $s13'$: *validTransO* (*lastt* $s1\ tr1, s1'$)
apply (*metis* $v3'$ *Opt.validFromS-def* *Opt.validS-append1* *Nil-is-append-conv* *hd-append2*)
by (*metis* *Opt.validFromS-def* *Opt.validS-validTrans* *lastt-def* *list.sel(1)* *not-Cons-self2*
snoc-eq-iff-butlast *tr1NE* $v3'$)

have *ninter3*: *never isIntO tr1* **and** *nis1'*: \neg *isIntO s1'*
using *ninter3'* **by** *auto*
obtain *ww ww1 ww2 trv1 trv2 statOO* **where** χ^3 : $\chi^3 \Delta ww w1 w2 ww1 ww2 s1$
tr1 s2 statA sv1 trv1 sv2 trv2 statOO
using *unwindCond-ex- χ^3 [OF unwind Δ r v3 nis1 nis2 ninter3 sec]* **by** *auto*

have *trv12NE*: *trv1* $\neq []$ *trv2* $\neq []$ **using** χ^3 **unfolding** χ^3 -*def* **by** *auto*

define *ss1 ssv1 ssv2* **where** *ss1*: *ss1* \equiv *lastt s1 tr1*
and *ssv1*: *ssv1* \equiv *lastt sv1 trv1* **and** *ssv2*: *ssv2* \equiv *lastt sv2 trv2*

have *ss1l*: *ss1* = *last tr1* **by** (*simp add: lastt-def ss1 tr1NE*)
have *tr1l*: *tr1* = *butlast tr1 @ [ss1]* **by** (*simp add: ss1 tr1NE*)
have *ssv1l*: *ssv1* = *last trv1* **using** χ^3 **unfolding** χ^3 -*def* **by** (*metis lastt-def ssv1*)
have *trv1l*: *trv1* = *butlast trv1 @ [ssv1]* **by** (*simp add: ssv1l trv12NE(1)*)
have *ssv2l*: *ssv2* = *last trv2* **using** χ^3 **unfolding** χ^3 -*def* **by** (*metis lastt-def ssv2*)
have *trv2l*: *trv2* = *butlast trv2 @ [ssv2]* **by** (*simp add: ssv2l trv12NE(2)*)

have *iss1[simp]*: *isSecO ss1* **using** *sec(2)* **unfolding** *ss1* **by** *auto*
have *issv1[simp]*: *isSecV ssv1* **and** *gissv13[simp]*: *getSecO ss1* = *getSecV ssv1*
using χ^3 **unfolding** χ^3 -*def ssv1 ss1* **by** *auto*

have *niss1*: \neg *isIntO ss1*
by (*metis list-all-append list-all-simps(1) ninter3 tr1l*)

have *rss1*: *reachO ss1* **and** *rssv12*: *reachV ssv1 reachV ssv2*
using *Opt.reach-validFromS-reach r ss1 tr1NE v3* **apply** *blast*
apply (*metis Van.reach-validFromS-reach χ^3 -def χ^3 r(3) ssv1l*)
by (*metis Van.reach-validFromS-reach χ^3 -def χ^3 r(4) ssv2l*)

have Δ : $\Delta ww ww1 ww2 ss1 s2 statA ssv1 ssv2 statOO$
using χ^3 **unfolding** χ^3 -*def ss1[symmetric] ssv1[symmetric] ssv2[symmetric]* **by** *auto*

have *s13'*: *validTransO (ss1, s1')*
by (*simp add: s13' ss1*)

note *vs13* = *s13'[unfolded ss1[symmetric]]*
obtain *w1' w2' trv1' trv2' statO'* **where**
 χ^{3a} : $\chi^{3a} \Delta \infty ww1 ww2 w1' w2' ss1 s1' s2 statA ssv1 trv1' ssv2 trv2' statO'$
using *unwindCond-ex- χ^{3a} -getSec[OF unwind Δ rss1 r(2) rssv12 s13' niss1 iss1*
issv1 gissv13]
by *blast*

have *trv12'NE*: *trv1'* $\neq []$ *trv2'* $\neq []$ **using** χ^{3a} **unfolding** χ^{3a} -*def* **by** *auto*

have [simp]: $\text{Van.O (butlast } trv1 \text{ @ } trv1') = \text{Van.O } trv1 \text{ @ Van.O } trv1'$
using $trv12'NE$ **unfolding** $\chi^3\text{-def}$ Van.O.map-filter Opt.O.map-filter **apply**(subst butlast-append) **by** simp

have [simp]: $\text{Van.O (butlast } trv2 \text{ @ } trv2') = \text{Van.O } trv2 \text{ @ Van.O } trv2'$
using $trv12'NE$ **unfolding** $\chi^3\text{-def}$ Van.O.map-filter Opt.O.map-filter **apply**(subst butlast-append) **by** simp

have $\text{Van.A } trv1' = \text{Van.A } trv2'$ **using** χ^3a **unfolding** $\chi^3a\text{-def}$ **by** auto
moreover **have** $\text{length (Van.O } trv1') = \text{length (Van.A } trv1') \wedge \text{length (Van.O } trv2') = \text{length (Van.A } trv2')$
unfolding Van.A.map-filter Van.O.map-filter **by** auto
ultimately **have** $\text{length (Van.O } trv1') = \text{length (Van.O } trv2')$ **by** auto
hence [simp]: $\text{Van.O } trv1 \text{ @ Van.O } trv1' = \text{Van.O } trv2 \text{ @ Van.O } trv2' \longleftrightarrow$
 $\text{Van.O } trv1 = \text{Van.O } trv2 \wedge \text{Van.O } trv1' = \text{Van.O } trv2'$ **by** auto

show ?thesis
apply(rule exI[of - $w1$]) **apply**(rule exI[of - $w2$])
apply(rule exI[of - butlast $trv1$ @ $trv1'$]) **apply**(rule exI[of - butlast $trv2$ @ $trv2'$])
apply(rule exI[of - statO])
unfolding $\chi^3b\text{-def}$ **apply**(intro conjI)
subgoal **using** χ^3 χ^3a **unfolding** $\chi^3\text{-def}$ $\chi^3a\text{-def}$ **by** auto
subgoal **using** χ^3 χ^3a **unfolding** $\chi^3\text{-def}$ $\chi^3a\text{-def}$ **by** auto
subgoal **using** χ^3 χ^3a **unfolding** $\chi^3\text{-def}$ $\chi^3a\text{-def}$
by simp (metis Simple-Transition-System.fromS-eq-Nil Simple-Transition-System.toS-fromS-nonSingl Van.toS-Nil diff-add-inverse2 linorder-not-less order-le-less-trans trans-less-add2 zero-less-diff)

subgoal **using** χ^3 χ^3a **unfolding** $\chi^3\text{-def}$ $\chi^3a\text{-def}$ ssv1
using $\text{Van.validFromS-append}$ **by** auto
subgoal **using** χ^3 χ^3a **unfolding** $\chi^3\text{-def}$ $\chi^3a\text{-def}$ ssv2
using $\text{Van.validFromS-append}$ **by** auto
subgoal **using** χ^3 χ^3a **unfolding** $\chi^3\text{-def}$ $\chi^3a\text{-def}$ **unfolding** Van.S.map-filter
 Opt.S.map-filter
apply(subst $tr1l$) **apply**(subst butlast-append)
by simp (metis Opt.S.map-filter $\text{Opt.S.eq-Nil-iff}(2)$ Van.S.map-filter $\text{Van.S.eq-Nil-iff}(2)$
 $\text{sec}(1)$)
subgoal **using** χ^3 χ^3a **unfolding** $\chi^3\text{-def}$ $\chi^3a\text{-def}$
by (simp add: butlast-append)
subgoal **using** χ^3 χ^3a **unfolding** $\chi^3\text{-def}$ $\chi^3a\text{-def}$ Van.A.map-filter Opt.A.map-filter

apply(subst $trv1l$) **apply**(subst $trv2l$) **by** (simp add: butlast-append)
subgoal **using** χ^3a $trv12'NE$ $tr1NE$ **unfolding** $\chi^3a\text{-def}$ lastt-def **by** simp .
qed

lemma unwindCond-ex- $\chi^3b\text{-aux}2$:
assumes unwind: unwindCond Δ
and Δ : $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$
and r : reachO $s1$ reachO $s2$ reachV $sv1$ reachV $sv2$

and $v3'$: $Opt.validFromS\ s1\ (tr1\ @\ [s1',s1''])$
and $nis1$: $\neg\ isIntO\ s1$ **and** $nis2$: $\neg\ isIntO\ s2$
and $ninter3'$: $never\ isIntO\ (tr1\ @\ [s1',s1''])$
and sec : $never\ isSecO\ tr1\ isSecO\ s1'$
shows $\exists\ w1'\ w2'\ trv1\ trv2\ statOO.\ \chi3b\ \Delta\ \infty\ w1\ w2\ w1'\ w2'\ s1\ (tr1\ @\ [s1',s1''])$
 $s2\ statA\ sv1\ trv1\ sv2\ trv2\ statOO$
proof –
have 0 : $lastt\ s1\ (tr1\ \#\#\ s1') = s1'$
unfolding $lastt-def$ **by** $auto$
show $?thesis$
using $unwindCond-ex-\chi3b-aux[OF\ unwind\ \Delta\ r,\ of\ tr1\ \#\#\ s1',\ unfolded\ 0,\ simplified]$
using $assms$ **by** $auto$
qed

definition $\chi3'\ \Delta\ w1\ w2\ w1'\ w2'\ s1\ tr1\ s1'\ s1''\ s2\ statAA\ sv1\ trv1\ sv1''\ sv2\ trv2\ sv2''\ statOO \equiv$
 $Van.validFromS\ sv1\ (trv1\ \#\#\ sv1'') \wedge Van.validFromS\ sv2\ (trv2\ \#\#\ sv2'') \wedge$
 $Van.S\ (trv1\ \#\#\ sv1'') = Opt.S\ ((tr1\ \#\#\ s1')\ \#\#\ s1'') \wedge never\ isSecV\ trv2 \wedge$
 $Van.A\ (trv1\ \#\#\ sv1'') = Van.A\ (trv2\ \#\#\ sv2'') \wedge$
 $trv1 \neq [] \wedge (trv2 \neq [] \vee w2' < w2) \wedge$
 $\Delta \infty\ w1'\ w2'\ s1''\ s2\ statAA\ sv1''\ sv2''\ statOO$

proposition $unwindCond-ex-\chi3'$:
assumes $unwind$: $unwindCond\ \Delta$
and Δ : $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ **and**
 r : $reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$
and $v3'$: $Opt.validFromS\ s1\ ((tr1\ \#\#\ s1')\ \#\#\ s1'')$
and $nis1$: $\neg\ isIntO\ s1$ **and** $nis2$: $\neg\ isIntO\ s2$
and $ninter3'$: $never\ isIntO\ ((tr1\ \#\#\ s1')\ \#\#\ s1'')$
and sec : $never\ isSecO\ tr1\ isSecO\ s1'$
shows $\exists\ w1'\ w2'\ trv1\ sv1''\ trv2\ sv2''\ statOO.\ \chi3'\ \Delta\ w1\ w2\ w1'\ w2'\ s1\ tr1\ s1'\ s1''\ s2\ statA\ sv1\ trv1\ sv1''\ sv2\ trv2\ sv2''\ statOO$
using $unwindCond-ex-\chi3b-aux2[unfolded\ \varphi-def,\ unfolded\ lastt-snoc\ lastt-snoc2\ append-snoc2,\ OF\ assms]$
unfolding $\chi3b-def$ **apply**($elim\ exE$) **subgoal for** $w1'\ w2'\ trv1\ trv2\ statOO$
apply($cases\ trv1\ rule:\ rev-cases$)
subgoal by $auto$
subgoal for $trv1'\ sv1''$ **apply**($cases\ trv2\ rule:\ rev-cases$)
subgoal by $auto$
subgoal for $trv2'\ sv2''$ **unfolding** $\chi3'-def$
apply($rule\ exI[of - w1']$) **apply**($rule\ exI[of - w2']$)
apply($rule\ exI[of - trv1']$) **apply**($rule\ exI[of - sv1'']$)
apply($rule\ exI[of - trv2']$) **apply**($rule\ exI[of - sv2'']$)
apply($rule\ exI[of - statOO]$)
by $simp\ (metis\ Opt.S.Nil-iff\ Opt.S.eq-Nil-iff(1)\ Van.S.simps(4)\ append-snoc2)$

list-all-append sec(2)
self-append-conv2 snoc-eq-iff-butlast) . . .

definition $\omega 3 \Delta w1 w2 w1' w2' s1 s1' s2 statAA sv1 trv1 sv1' sv2 trv2 sv2'$
 $statOO \equiv$
Van.validFromS sv1 (trv1 ## sv1') \wedge Van.validFromS sv2 (trv2 ## sv2') \wedge
never isSecV trv1 \wedge never isSecV trv2 \wedge
Van.A (trv1 ## sv1') = Van.A (trv2 ## sv2') \wedge
(trv1 \neq [] \vee w1' < w1) \wedge (trv2 \neq [] \vee w2' < w2) \wedge
 $\Delta \infty w1' w2' s1' s2 statAA sv1' sv2' statOO$

proposition *unwindCond-ex- $\omega 3$:*
assumes *unwind: unwindCond Δ*
and $\Delta: \Delta w w1 w2 s1 s2 statA sv1 sv2 statO$
and $r34: reachO s1 reachO s2$ **and** $r12: reachV sv1 reachV sv2$
and $v3: validTransO (s1, s1')$
and $nis1: \neg isIntO s1 \neg isIntO s1' \neg isSecO s1$
and $nis2: \neg isIntO s2$
shows $\exists w1' w2' trv1 sv1' trv2 sv2' statOO. \omega 3 \Delta w1 w2 w1' w2' s1 s1' s2 statA$
 $sv1 trv1 sv1' sv2 trv2 sv2' statOO$
using $\Delta r12$
proof(*induction w arbitrary: w1 w2 sv1 sv2 statO rule: less-induct*)
case (*less w w1 w2 sv1 sv2 statO*)
note $\Delta = \langle \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \rangle$
note $r12 = less.prem(2,3)$
note $r1 = r12(1)$ **note** $r2 = r12(2)$
note $r = r34 r12$

have $f34: finalO s1 = finalO s2 \wedge finalV sv1 = finalO s1 \wedge finalV sv2 = finalO$
 $s2$
using $\Delta unwind[unfolded unwindCond-def] r$ **by** *auto*

have *proact-match: ($\exists v < w. proact \Delta v w1 w2 s1 s2 statA sv1 sv2 statO$) \vee match*
 $\Delta w1 w2 s1 s2 statA sv1 sv2 statO$
using $\Delta unwind[unfolded unwindCond-def] r$ **by** *auto*
show *?case using proact-match proof safe*
fix v **assume** $v: v < w$
assume *proact $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$*
thus *?thesis unfolding proact-def proof safe*
assume $sv1: \neg isSecV sv1 \neg isIntV sv1$ **and** *move-1 $\Delta v w1 w2 s1 s2 statA$*
 $sv1 sv2 statO$
then obtain $sv1'$ **where** $0: validTransV (sv1, sv1')$ **and** $\Delta: \Delta v w1 w2 s1$
 $s2 statA sv1' sv2 statO$
unfolding *move-1-def by auto*
have $r1': reachV sv1'$ **using** $r1 0$ **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain $w1' w2' trv1 sv1'' trv2 sv2' statOO$ **where**
 $\omega 3: \omega 3 \Delta w1 w2 w1' w2' s1 s1' s2 statA sv1' trv1 sv1'' sv2 trv2 sv2' statOO$

```

    using less(1)[OF v Δ r1' r2] by auto
    show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
exI[of - sv1 # trv1]) apply(rule exI[of - sv1''])
    apply(rule exI[of - trv2]) apply(rule exI[of - sv2'])
    using ω3 0 sv1 unfolding ω3-def by auto
next
    assume sv2: ¬ isSecV sv2 ¬ isIntV sv2 and move-2 Δ v w1 w2 s1 s2 statA
sv1 sv2 statO
    then obtain sv2'
    where 0: validTransV (sv2,sv2')
    and Δ: Δ v w1 w2 s1 s2 statA sv1 sv2' statO
    unfolding move-2-def by auto
    have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
    obtain w1' w2' trv1 sv1'' trv2 sv2'' statOO where
    ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1'' sv2' trv2 sv2'' statOO

    using less(1)[OF v Δ r1 r2'] by auto
    show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - trv1]) apply(rule exI[of - sv1'])
    apply(rule exI[of - sv2 # trv2]) apply(rule exI[of - sv2''])
    using ω3 0 sv2 unfolding ω3-def by auto
next
    assume sv1: ¬ isSecV sv1 and sv2: ¬ isSecV sv2 and
    move-12 Δ v w1 w2 s1 s2 statA sv1 sv2 statO and
    sv12: Van.eqAct sv1 sv2
    then obtain sv1' sv2' statO'
    where statO': statO' = sstatO' statO sv1 sv2
    and 0: validTransV (sv1,sv1') validTransV (sv2,sv2')
    and Δ: Δ v w1 w2 s1 s2 statA sv1' sv2' statO'
    unfolding move-12-def by auto
    have r1': reachV sv1' and r2': reachV sv2' using r1 r2 0
    by (metis Van.reach.Step fst-conv snd-conv)+
    obtain w1' w2' trv1 sv1'' trv2 sv2'' statOO where
    ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1' trv1 sv1'' sv2' trv2 sv2'' statOO

    using less(1)[OF v Δ r1' r2'] by auto
    show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv1''])
    apply(rule exI[of - sv2 # trv2]) apply(rule exI[of - sv2''])
    using ω3 0 sv1 sv2 sv12 unfolding ω3-def statO' by (auto simp: Van.eqAct-def)
qed
next
    assume m: match Δ w1 w2 s1 s2 statA sv1 sv2 statO
    have m: match1 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
match-def by auto
    have (∃ w1' w2'. w1' < w1 ∧ w2' < w2 ∧ ¬ isSecO s1 ∧ Δ ∞ w1' w2' s1' s2
statA sv1 sv2 statO) ∨
    (∃ w2' < w2. eqSec sv1 s1 ∧ ¬ isIntV sv1 ∧ match1-1 Δ ∞ w2' s1 s1' s2
statA sv1 sv2 statO) ∨

```

(*eqSec sv1 s1* \wedge \neg *isSecV sv2* \wedge *Van.eqAct sv1 sv2* \wedge *match1-12* Δ ∞ ∞
s1 s1' s2 statA sv1 sv2 statO)

using *m v3 nis1 unfolding match1-def by auto*

thus *?thesis*

apply(*elim disjE exE*)

subgoal for *w1' w2'*

apply(*rule exI[of - w1']*) **apply**(*rule exI[of - w2']*)

apply(*rule exI[of - []]*) **apply**(*rule exI[of - sv1]*)

apply(*rule exI[of - []]*) **apply**(*rule exI[of - sv2]*)

apply(*rule exI[of - statO]*) **unfolding** $\omega 3$ -*def*

by auto

subgoal for *w2'*

apply(*rule exI[of - ∞]*) **apply**(*rule exI[of - w2']*)

unfolding *match1-1-def* **apply**(*elim conjE exE*) **subgoal for** *sv1'*

apply(*rule exI[of - [sv1]]*) **apply**(*rule exI[of - sv1']*)

apply(*rule exI[of - []]*) **apply**(*rule exI[of - sv2]*)

apply(*rule exI[of - statO]*)

unfolding $\omega 3$ -*def* **using** *nis1(3)* **by** (*auto simp: eqSec-def*) .

subgoal

apply(*rule exI[of - ∞]*) **apply**(*rule exI[of - ∞]*)

unfolding *match1-12-def* **apply**(*elim conjE exE*) **subgoal for** *sv1' sv2'*

apply(*rule exI[of - [sv1]]*) **apply**(*rule exI[of - sv1']*)

apply(*rule exI[of - [sv2]]*) **apply**(*rule exI[of - sv2']*)

apply(*rule exI[of - sstatO' statO sv1 sv2]*)

unfolding $\omega 3$ -*def* **using** *nis1(3)* **apply** (*auto simp: eqSec-def*

sstatA'-def sstatO'-def lastt-def Van.eqAct-def) . . .

qed

qed

definition $\chi 4$ Δ *w w1 (w2::enat) w1' w2' s1 s2 tr2 statAA sv1 trv1 sv2 trv2*
statOO \equiv

trv1 \neq [] \wedge *trv2* \neq [] \wedge (*length trv1* $>$ *Suc 0* \vee *w1' \leq w1*) \wedge

Van.validFromS sv1 trv1 \wedge *Van.validFromS sv2 trv2* \wedge

never isSecV (butlast trv1) \wedge

never isSecV (butlast trv2) \wedge

isSecV (lastt sv2 trv2) \wedge *getSecV (lastt sv2 trv2) = getSecO (lastt s2 tr2)* \wedge

Van.A trv1 = Van.A trv2 \wedge

Δ *w w1' w2' s1 (lastt s2 tr2) statAA (lastt sv1 trv1) (lastt sv2 trv2) statOO*

lemma $\chi 4$ -*final*:

assumes *unw: unwindCond* Δ

and *r: reachO s1 reachO s2 reachV sv1 reachV sv2*

and *vtr2: Opt.validFromS s2 tr2*

and $\chi 4$: $\chi 4$ Δ *w w1 w2 w1' w2' s1 s2 tr2 statAA sv1 trv1 sv2 trv2 statOO*

shows (*finalV (lastt sv1 trv1)* \longleftrightarrow *finalO s1*) \wedge (*finalV (lastt sv2 trv2)* \longleftrightarrow *finalO*)

(*lastt s2 tr2*)
proof –
have *rsv12*: *Van.validFromS sv1 trv1* \longrightarrow *reachV (lastt sv1 trv1)*
 Van.validFromS sv2 trv2 \longrightarrow *reachV (lastt sv2 trv2)* **using** *r*
by (*simp add: Van.reach-validFromS-reach lastt-def*)
have *rs2*: *Opt.validFromS s2 tr2* \longrightarrow *reachO (lastt s2 tr2)*
using *r*
by (*simp add: Opt.reach-validFromS-reach lastt-def*)
show *?thesis* **using** χ_4 [*unfolded* χ_4 -*def*] *rsv12 rs2* **using** *unw[unfolded unwind-Cond-def, rule-format,*
 of s1 lastt s2 tr2 lastt sv1 trv1 lastt sv2 trv2 w w1' w2' statAA statOO]
using *vtr2* \langle *reachO s1* \rangle **by** *auto*
qed

lemma χ_4 -*completedFrom*: *unwindCond* $\Delta \implies$
reachO s1 \implies *reachO s2* \implies *reachV sv1* \implies *reachV sv2* \implies
Opt.validFromS s2 tr2 \implies *completedFromO s2 tr2* \implies
 χ_4 Δ *w w1 w2 w1' w2' s1 s2 tr2 statAA sv1 trv1 sv2 trv2 statOO*
 \implies *completedFromV sv1 trv1* \wedge *completedFromV sv2 trv2*
by (*metis Van.final-not-isSec* χ_4 -*def* χ_4 -*final completedFromO-lastt*)

proposition *unwindCond-ex- χ_4* :
assumes *unwind*: *unwindCond* Δ
and Δ : Δ *w w1 w2 s1 s2 statA sv1 sv2 statO*
and *r*: *reachO s1 reachO s2 reachV sv1 reachV sv2*
and *vtr2*: *Opt.validFromS s2 tr2*
and *nis2*: \neg *isIntO s1* **and** *nis2*: \neg *isIntO s2*
and *inter4*: *never isIntO tr2*
and *sec*: *never isSecO (butlast tr2) isSecO (lastt s2 tr2)*
shows \exists *w' w1' w2' trv1 trv2 statOO*. χ_4 Δ *w' w1 w2 w1' w2' s1 s2 tr2 statA sv1*
trv1 sv2 trv2 statOO
using *assms(2-)*
proof(*induction length tr2 w*
 arbitrary: w1 w2 s1 s2 statA sv1 sv2 statO tr2 rule: less2-induct')
case (*less w tr2 w1 w2 s1 s2 statA sv1 sv2 statO*)
note *vtr2* = *less(8)*
note Δ = \langle Δ *w w1 w2 s1 s2 statA sv1 sv2 statO* \rangle
note *nis1* = *less(9)* **note** *nis2* = *less(10)*
note *inter4* = *less(11)*
note *sec4* = *less(12,13)*
note *r34* = *less.prem(2,3)* **note** *r12* = *less.prem(4,5)*
note *r* = *r34 r12*
note *r3* = *r34(1)* **note** *r4* = *r34(2)* **note** *r1* = *r12(1)* **note** *r2* = *r12(2)*

have *i34*: *statA = Eq* \longrightarrow *isIntO s1 = isIntO s2*
and *f34*: *finalO s1 = finalO s2* \wedge *finalV sv1 = finalO s1* \wedge *finalV sv2 = finalO*
s2
using Δ *unwind[unfolded unwindCond-def]* *r* **by** *auto*

```

have proact-match: ( $\exists v < w. \text{proact } \Delta v w1 w2 s1 s2 \text{ statA } sv1 sv2 \text{ statO}$ )  $\vee$  match
 $\Delta w1 w2 s1 s2 \text{ statA } sv1 sv2 \text{ statO}$ 
using  $\Delta \text{unwind}[\text{unfolded unwindCond-def}] r$  by auto
show ?case using proact-match proof safe
  fix v assume v:  $v < w$ 
  assume proact  $\Delta v w1 w2 s1 s2 \text{ statA } sv1 sv2 \text{ statO}$ 
  thus ?thesis unfolding proact-def proof safe
  assume sv1:  $\neg \text{isSecV } sv1 \neg \text{isIntV } sv1$  and move-1  $\Delta v w1 w2 s1 s2 \text{ statA}$ 
 $sv1 sv2 \text{ statO}$ 
  then obtain sv1'
  where 0: validTransV (sv1, sv1')
  and  $\Delta$ :  $\Delta v w1 w2 s1 s2 \text{ statA } sv1' sv2 \text{ statO}$ 
  unfolding move-1-def by auto
  have r1': reachV sv1' using r1 0 by (metis Van.reach.Step fst-conv snd-conv)
  obtain w' w1' w2' trv1 trv2 statOO where  $\chi_4$ :  $\chi_4 \Delta w' w1 w2 w1' w2' s1$ 
 $s2 \text{ tr2 } \text{statA } sv1' \text{ trv1 } sv2 \text{ trv2 } \text{statOO}$ 
  using less(2)[OF v, of tr2 w1 w2 s1 s2 statA sv1' sv2 statO,
simplified, OF  $\Delta r3_4 r1' r2 \text{ vtr2 nis1 nis2 inter}_4 \text{ sec}_4]$  by auto
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
  using  $\chi_4 0 \text{ sv1}$  unfolding  $\chi_4$ -def by auto
  next
  assume sv2:  $\neg \text{isSecV } sv2 \neg \text{isIntV } sv2$  and move-2  $\Delta v w1 w2 s1 s2 \text{ statA}$ 
 $sv1 sv2 \text{ statO}$ 
  then obtain sv2'
  where 0: validTransV (sv2, sv2')
  and  $\Delta$ :  $\Delta v w1 w2 s1 s2 \text{ statA } sv1 sv2' \text{ statO}$ 
  unfolding move-2-def by auto
  have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
  obtain w1' w2' w' trv1 trv2 statOO where  $\chi_4$ :  $\chi_4 \Delta w' w1 w2 w1' w2' s1$ 
 $s2 \text{ tr2 } \text{statA } sv1 \text{ trv1 } sv2' \text{ trv2 } \text{statOO}$ 
  using less(2)[OF v, of tr2 w1 w2 s1 s2 statA sv1 sv2' statO,
simplified, OF  $\Delta r3_4 r1 r2' \text{ vtr2 nis1 nis2 inter}_4 \text{ sec}_4]$  by auto
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
  using  $\chi_4 0 \text{ sv2}$  unfolding  $\chi_4$ -def by auto
  next
  assume sv12:  $\neg \text{isSecV } sv1 \neg \text{isSecV } sv2$  Van.eqAct sv1 sv2
  and move-12  $\Delta v w1 w2 s1 s2 \text{ statA } sv1 sv2 \text{ statO}$ 
  then obtain sv1' sv2' statO'
  where 0: statO' = sstatO' statO sv1 sv2
validTransV (sv1, sv1')  $\neg \text{isSecV } sv1$ 
validTransV (sv2, sv2')  $\neg \text{isSecV } sv2$ 
Van.eqAct sv1 sv2
  and  $\Delta$ :  $\Delta v w1 w2 s1 s2 \text{ statA } sv1' sv2' \text{ statO}'$ 
  unfolding move-12-def by auto
  have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
fst-conv snd-conv) $+$ 
  obtain w' w1' w2' trv1 trv2 statOO where  $\chi_4$ :  $\chi_4 \Delta w' w1 w2 w1' w2' s1$ 

```

```

s2 tr2 statA sv1' trv1 sv2' trv2 statOO
  using less(2)[OF v, of tr2 w1 w2 s1 s2 statA sv1' sv2' statO',
    simplified, OF Δ r34 r12' vtr2 nis1 nis2 inter4 sec4] by auto
  show ?thesis apply(rule exI[of - w]) apply(rule exI[of - w1]) apply(rule
exI[of - w2]) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
  apply(rule exI[of - statOO])
  using χ4 0 sv12 unfolding χ4-def sstatO'-def
  by (auto simp: Van.eqAct-def)
qed
next
assume m: match Δ w1 w2 s1 s2 statA sv1 sv2 statO
define statA' where statA': statA' = sstatA' statA s1 s2
show ?thesis
proof (cases length tr2 ≤ Suc 0)
  case True
  hence tr2e: tr2 = [] ∨ tr2 = [s2]
  by (metis Opt.validFromS-def Suc-length-conv le-Suc-eq le-zero-eq length-0-conv
list.sel(1) vtr2)
  hence Opt.A tr2 = [] by (simp add: True)
  have is2: isSecO s2
  by (metis last.simps lastt-def sec4(2) tr2e)
  hence ¬ finalO s2 using Opt.final-not-isSec by blast
  then obtain s2' where s24': validTransO (s2, s2') unfolding Opt.final-def
by auto
  hence isv2: isSecV sv2 ∧ getSecV sv2 = getSecO s2 using m is2 nis2
  unfolding match-def match2-def eqSec-def by auto
  show ?thesis using tr2e isv2 apply-
  apply(rule exI[of - w]) apply(rule exI[of - w1]) apply(rule exI[of - w2])
  apply(rule exI[of - [sv1]], rule exI[of - [sv2]], rule exI[of - statO])
  using tr2e
  using f34 Δ by (clarsimp simp: χ4-def lastt-def)
next
case False
then obtain s24 tr2' where tr2: tr2 = s24 # tr2' and tr2'NE: tr2' ≠ []
  by (cases tr2, auto)
have s24[simp]: s24 = s2 using ⟨Opt.validFromS s2 tr2⟩
  by (simp add: Opt.validFromS-Cons-iff tr2)
obtain s2' where
  trn4: validTransO (s2, s2') and
  tr2': Opt.validFromS s2' tr2' using ⟨Opt.validFromS s2 tr2⟩
unfolding tr2 s24 by (metis tr2'NE Simple-Transition-System.validFromS-Cons-iff)
have r4': reachO s2' using r4 trn4 by (metis Opt.reach.Step fst-conv snd-conv)
have f4: ¬ finalO s2 using Opt.final-def trn4 by blast
hence f3: ¬ finalO s1 using f34 by blast
have nev4': never isIntO tr2'
using inter4 tr2 tr2'NE by auto
have isAO4: ¬ isIntO s2 by (simp add: nis2)
have O44': Opt.O tr2 = Opt.O tr2' Opt.A tr2 = Opt.A tr2'
using isAO4 unfolding tr2 by auto

```

have m : $\text{match2 } \Delta w1 w2 s1 s2 \text{ statA } sv1 sv2 \text{ statO}$ **using** m **unfolding**
match-def **by** *auto*
have $(\exists w1' < w1. \exists w2' < w2. \neg \text{isSecO } s2 \wedge \Delta \infty w1' w2' s1 s2' \text{ statA } sv1$
 $sv2 \text{ statO}) \vee$
 $(\exists w1' < w1. \text{eqSec } sv2 s2 \wedge \neg \text{isIntV } sv2 \wedge \text{match2-1 } \Delta w1' \infty s1 s2 s2'$
 $\text{statA } sv1 sv2 \text{ statO}) \vee$
 $(\text{eqSec } sv2 s2 \wedge \neg \text{isSecV } sv1 \wedge \text{Van.eqAct } sv1 sv2 \wedge \text{match2-12 } \Delta \infty \infty$
 $s1 s2 s2' \text{ statA } sv1 sv2 \text{ statO})$
using m $\text{isAO4 } trn4$ **unfolding** *match2-def* **by** *auto*
thus *?thesis*
proof *safe*
fix $w1'' w2''$ **assume** $w12'$: $w1'' < w1 w2'' < w2$
assume $\neg \text{isSecO } s2$ **and** Δ : $\Delta \infty w1'' w2'' s1 s2' \text{ statA } sv1 sv2 \text{ statO}$
hence $S4$: $\text{Opt.S } tr2' = \text{Opt.S } tr2$ **unfolding** $tr2$ **by** *auto*
obtain $w' w1' w2' trv1 trv2 \text{ statOO}$ **where** $\chi4$: $\chi4 \Delta w' w1'' w2'' w1' w2'$
 $s1 s2' tr2' \text{ statA } sv1 trv1 sv2 trv2 \text{ statOO}$
using $\text{less}(1)[\text{of } tr2', \text{OF} - \Delta r3 r4' r12]$ **unfolding** $tr2$
by *simp* (*metis* $\text{Opt.S.eq-Nil-iff}(2)$ $S4$ *Simple-Transition-System.validFromS-def*
last.simps lastt-def
 $\text{list.discI list-all-hd nev4' nis1 sec4}(1) \text{ sec4}(2) tr2 tr2' tr2'NE$)
show *?thesis* **apply**(*rule* $\text{exI}[\text{of} - w']$) **apply**(*rule* $\text{exI}[\text{of} - w1']$) **apply**(*rule*
 $\text{exI}[\text{of} - w2']$) **apply**(*rule* $\text{exI}[\text{of} - trv1]$) **apply**(*rule* $\text{exI}[\text{of} - trv2]$)
using $\chi4$ $O44'$ **unfolding** $\chi4\text{-def } tr2$ *Van.completedFrom-def*
using *Van.validFromS-Cons* $tr2'NE$ $tr2' trn4 \text{isAO4 } w12'$ **by** *auto*
next
fix $w1''$ **assume** $w1'$: $w1'' < w1$
assume $trn24$: $\text{eqSec } sv2 s2$ **and**
 Atrn2 : $\neg \text{isIntV } sv2$ **and** $\text{match2-1 } \Delta w1'' \infty s1 s2 s2' \text{ statA } sv1 sv2 \text{ statO}$
then obtain $sv2'$ **where**
 $trn2$: $\text{validTransV } (sv2, sv2')$ **and**
 Δ : $\Delta \infty w1'' \infty s1 s2' \text{ statA } sv1 sv2' \text{ statO}$
unfolding *match2-1-def* **by** *auto*
have $r2'$: $\text{reachV } sv2'$ **using** $r2$ $trn2$ **by** (*metis* *Van.reach.Step fst-conv*
snd-conv)
obtain $w' w1' w2' trv1 trv2 \text{ statOO}$ **where** $\chi4$: $\chi4 \Delta w' w1'' \infty w1' w2'$
 $s1 s2' tr2' \text{ statA } sv1 trv1 sv2' trv2 \text{ statOO}$
using $\text{less}(1)[\text{of } tr2', \text{OF} - \Delta r3 r4' r1 r2', \text{unfolded } O44', \text{simplified}]$
using $\text{less.premis } tr2' f3 f4 tr2'NE$ $trn4$ $O44'(1)$
unfolding $tr2$
by *simp* (*metis* *Opt.validFromS-def list-all-hd*)
show *?thesis* **apply**(*rule* $\text{exI}[\text{of} - w']$) **apply**(*rule* $\text{exI}[\text{of} - w1']$) **apply**(*rule*
 $\text{exI}[\text{of} - w2']$) **apply**(*rule* $\text{exI}[\text{of} - trv1]$) **apply**(*rule* $\text{exI}[\text{of} - sv2 \# trv2]$)
using $\chi4$ $O44'$ **unfolding** $\chi4\text{-def } tr2$ *Van.completedFrom-def*
using *Van.validFromS-Cons* $trn2$ $tr2'NE$ $tr2' trn4$
using $\text{isAO4 } \text{Atrn2 } \text{eqSec-S-Cons } trn24$ $w1'$
by *simp* (*metis* $\text{Opt.S.Nil-iff } \text{Opt.S.eq-Nil-iff}(1)$ *eqSec-def nless-le order-le-less-trans*
 $s24 \text{ sec4}(1) tr2$)
next
assume $sv1$: $\neg \text{isSecV } sv1$ **and** $trn24$: $\text{eqSec } sv2 s2$ **and**

Atrn12: *Van.eqAct sv1 sv2 and match2-12* $\Delta \infty \infty s1 s2 s2' statA sv1 sv2$
statO
then obtain *sv1' sv2' statO'* **where**
statO': statO' = sstatO' statO sv1 sv2 and
trn1: validTransV (sv1,sv1') and
trn2: validTransV (sv2,sv2') and
 $\Delta: \Delta \infty \infty \infty s1 s2' statA sv1' sv2' statO'$
unfolding *match2-12-def* **by** *auto*
have *r12': reachV sv1' reachV sv2'*
using *r1 trn1 r2 trn2* **by** *(metis Van.reach.Step fst-conv snd-conv)+*
obtain *w' w1' w2' trv1 trv2 statOO* **where** $\chi_4: \chi_4 \Delta w' \infty \infty w1' w2' s1$
s2' tr2' statA sv1' trv1 sv2' trv2 statOO
using *less(1)[of tr2', OF - Δ r3 r4' r12', unfolded O44', simplified]*
using *less.premis tr2' f3 f4 tr2'NE trn4 O44'(1)* **unfolding** *tr2 statO'*
sstatO'-def
by *simp (metis Simple-Transition-System.validFromS-def list-all-hd)*
show *?thesis* **apply**(*rule exI[of - w']*) **apply**(*rule exI[of - w1']*) **apply**(*rule*
exI[of - w2'])
apply(*rule exI[of - sv1 # trv1]*) **apply**(*rule exI[of - sv2 # trv2]*)
using χ_4 *O44' tr2'NE sv1*
using *Van.validFromS-Cons trn1 trn2*
using *isAO4 Atrn12 eqSec-S-Cons trn24 f3 f34 s24 tr2' trn4*
unfolding χ_4 -*def tr2 Van.completedFrom-def Van.eqAct-def*
using *Van.A.Cons-unfold eqSec-def sec4(1) tr2* **by** *auto*
qed
qed
qed
qed

definition χ_4a **where** $\chi_4a \Delta w w1 (w2::enat) w1' w2' s1 s2 s2' statAA sv1 trv1$
sv2 trv2 statOO \equiv
 $trv1 \neq [] \wedge trv2 \neq [] \wedge (length\ trv1 > Suc\ 0 \vee w1' < w1) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge Van.validFromS\ sv2\ trv2 \wedge$
 $never\ isSecV\ (butlast\ trv1) \wedge$
 $Van.S\ trv2 = [getSecO\ s2] \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $\Delta\ w\ w1'\ w2'\ s1\ s2'\ statAA\ (lastt\ sv1\ trv1)\ (lastt\ sv2\ trv2)\ statOO$

lemma *unwindCond-ex- χ_4a -getSec:*
assumes *unwind: unwindCond Δ*
and $\Delta: \Delta w w1 w2 s1 s2 statA sv1 sv2 statO$
and *r34: reachO s1 reachO s2 and r12: reachV sv1 reachV sv2*
and *v: validTransO (s2, s2')*
and *ii4: $\neg isIntO s2$*
and *is2: isSecO s2 and isv24: isSecV sv2 getSecO s2 = getSecV sv2*
shows $\exists w1' w2' trv1 trv2 statOO.$
 $\chi_4a \Delta \infty w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv2 trv2 statOO$
using Δ *r12 isv24*
proof(*induction w arbitrary: w1 w2 sv1 sv2 statO rule: less-induct*)

```

case (less w w1 w2 sv1 sv2 statO)
note Δ = ⟨Δ w w1 w2 s1 s2 statA sv1 sv2 statO⟩
note r12 = less.prem(2,3)
note r1 = r12(1) note r2 = r12(2)
note r = r34 r12
note isv24 = ⟨isSecV sv2⟩ ⟨getSecO s2 = getSecV sv2⟩

have f34: finalO s1 = finalO s2 ∧ finalV sv1 = finalO s1 ∧ finalV sv2 = finalO
s2
using Δ unwind[unfolded unwindCond-def] r by auto

have proact-match: (∃ v < w. proact Δ v w1 w2 s1 s2 statA sv1 sv2 statO) ∨ match
Δ w1 w2 s1 s2 statA sv1 sv2 statO
using Δ unwind[unfolded unwindCond-def] r by auto
show ?case using proact-match proof safe
fix v assume v: v < w
assume proact Δ v w1 w2 s1 s2 statA sv1 sv2 statO
thus ?thesis unfolding proact-def proof safe
assume sv1: ¬ isSecV sv1 ¬ isIntV sv1 and move-1 Δ v w1 w2 s1 s2 statA
sv1 sv2 statO
then obtain sv1'
where 0: validTransV (sv1,sv1')
and Δ: Δ v w1 w2 s1 s2 statA sv1' sv2 statO
unfolding move-1-def by auto
have r1': reachV sv1' using r1 0 by (metis Van.reach.Step fst-conv snd-conv)
obtain w1' w2' trv1 trv2 statOO where
χ4a: χ4a Δ ∞ w1 w2 w1' w2' s1 s2 s2' statA sv1' trv1 sv2 trv2 statOO
using less(1)[OF v Δ r1' r2 isv24] by auto
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
using χ4a 0 sv1 unfolding χ4a-def by auto
next
assume sv2: ¬ isSecV sv2 ¬ isIntV sv2 and move-2 Δ v w1 w2 s1 s2 statA
sv1 sv2 statO
hence False using isv24 by blast
thus ?thesis by auto
next
assume sv12: ¬ isSecV sv1 ¬ isSecV sv2 Van.eqAct sv1 sv2
and move-12 Δ v w1 w2 s1 s2 statA sv1 sv2 statO
hence False using isv24 by blast
thus ?thesis by auto
qed
next
assume m: match Δ w1 w2 s1 s2 statA sv1 sv2 statO
have m: match2 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
match-def by auto
have (∃ w1' w2'. w1' < w1 ∧ w2' < w2 ∧ ¬ isSecO s2 ∧ Δ ∞ w1' w2' s1 s2'
statA sv1 sv2 statO) ∨
(∃ w1' < w1. eqSec sv2 s2 ∧ ¬ isIntV sv2 ∧ match2-1 Δ w1' ∞ s1 s2 s2')

```

$statA\ sv1\ sv2\ statO) \vee$
 $(eqSec\ sv2\ s2 \wedge \neg isSecV\ sv1 \wedge Van.eqAct\ sv1\ sv2 \wedge match2-12\ \Delta\ \infty\ \infty$
 $s1\ s2\ s2'\ statA\ sv1\ sv2\ statO)$
using $m\ v\ ii4$ **unfolding** $match2-def$ **by** $auto$

thus $?thesis$
apply $(elim\ disjE\ exE)$
subgoal for $w1'\ w2'$ **using** $is2$ **by** $auto$
subgoal for $w1'$ **apply** $(rule\ exI[of - w1'])$ **apply** $(rule\ exI[of - \infty])$
unfolding $match2-1-def$ **apply** $(elim\ conjE\ exE)$ **subgoal for** $sv2'$
apply $(rule\ exI[of - [sv1]])$ **apply** $(rule\ exI[of - [sv2,sv2']])$
apply $(rule\ exI[of - statO])$
using $is2\ isv24$ **unfolding** $\chi4a-def$
by $(auto\ simp : sstatA'-def\ lastt-def) .$
subgoal **apply** $(rule\ exI[of - \infty])$ **apply** $(rule\ exI[of - \infty])$
unfolding $match2-12-def$ **apply** $(elim\ conjE\ exE)$ **subgoal for** $sv1'\ sv2'$
apply $(rule\ exI[of - [sv1,sv1']])$ **apply** $(rule\ exI[of - [sv2,sv2']])$
apply $(rule\ exI[of - sstatO'\ statO\ sv1\ sv2])$
using $is2\ isv24$ **unfolding** $\chi4a-def$
by $(auto\ simp : sstatA'-def\ sstatO'-def\ lastt-def\ Van.eqAct-def) . .$

qed

qed

definition $\chi4b\ \Delta\ w\ w1\ w2\ w1'\ (w2'::enat)\ s1\ s2\ tr2\ statAA\ sv1\ trv1\ sv2\ trv2$
 $statOO \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge (length\ trv1 > Suc\ 0 \vee w1' < w1) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge Van.validFromS\ sv2\ trv2 \wedge$
 $never\ isSecV\ (butlast\ trv1) \wedge$
 $Van.S\ trv2 = Opt.S\ tr2 \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $\Delta\ w\ w1'\ w2'\ s1\ (lastt\ s2\ tr2)\ statAA\ (lastt\ sv1\ trv1)\ (lastt\ sv2\ trv2)\ statOO$

lemma $unwindCond-ex-\chi4b-aux:$

assumes $unwind: unwindCond\ \Delta$

and $\Delta: \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and $r: reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$

and $tr2NE: tr2 \neq []$

and $v4': Opt.validFromS\ s2\ (tr2\ ##\ s2')$

and $nis1: \neg isIntO\ s1$ **and** $nis2: \neg isIntO\ s2$

and $ninter4': never\ isIntO\ (tr2\ ##\ s2')$

and $sec: never\ isSecO\ (butlast\ tr2)\ isSecO\ (lastt\ s2\ tr2)$

shows $\exists w1'\ w2'\ trv1\ trv2\ statOO. \chi4b\ \Delta\ \infty\ w1\ w2\ w1'\ w2'\ s1\ s2\ (tr2\ ##\ s2')$

$statA\ sv1\ trv1\ sv2\ trv2\ statOO$

proof –

have $v4: Opt.validFromS\ s2\ tr2$ **and** $s24': validTransO\ (lastt\ s2\ tr2, s2')$

apply $(metis\ v4'\ Opt.validFromS-def\ Opt.validS-append1\ Nil-is-append-conv\ hd-append2)$

by $(metis\ Opt.validFromS-def\ Opt.validS-validTrans\ append-is-Nil-conv\ lastt-def\ list.distinct(1)\ list.sel(1)\ tr2NE\ v4')$

have $ninter4$: *never isIntO tr2* **and** $nis2'$: $\neg isIntO s2'$
using $ninter4'$ **by** *auto*
obtain $ww\ ww1\ ww2\ trv1\ trv2\ statOO$ **where** $\chi4$: $\chi4\ \Delta\ ww\ w1\ w2\ ww1\ ww2\ s1\ s2\ tr2\ statA\ sv1\ trv1\ sv2\ trv2\ statOO$
using $unwindCond-ex-\chi4$ [$OF\ unwind\ \Delta\ r\ v4\ nis1\ nis2\ ninter4\ sec$]
by *auto*

have $trv12NE$: $trv1 \neq []\ trv2 \neq []$ **using** $\chi4$ **unfolding** $\chi4-def$ **by** *auto*

define $ss2\ ssv1\ ssv2$ **where** $ss2$: $ss2 \equiv lastt\ s2\ tr2$
and $ssv1$: $ssv1 \equiv lastt\ sv1\ trv1$ **and** $ssv2$: $ssv2 \equiv lastt\ sv2\ trv2$

have $ss2l$: $ss2 = last\ tr2$ **by** (*simp add: lastt-def ss2 tr2NE*)
have $tr2l$: $tr2 = butlast\ tr2\ @\ [ss2]$ **by** (*simp add: ss2l tr2NE*)
have $ssv1l$: $ssv1 = last\ trv1$ **using** $\chi4$ **unfolding** $\chi4-def$ **by** (*metis lastt-def ssv1*)
have $trv1l$: $trv1 = butlast\ trv1\ @\ [ssv1]$ **by** (*simp add: ssv1l trv12NE(1)*)
have $ssv2l$: $ssv2 = last\ trv2$ **using** $\chi4$ **unfolding** $\chi4-def$ **by** (*metis lastt-def ssv2*)
have $trv2l$: $trv2 = butlast\ trv2\ @\ [ssv2]$ **by** (*simp add: ssv2l trv12NE(2)*)

have $iss2[simp]$: *isSecO ss2* **using** $sec(2)$ **unfolding** $ss2$ **by** *auto*
have $issv2[simp]$: *isSecV ssv2* **and** $gissv24[simp]$: *getSecO ss2 = getSecV ssv2*
using $\chi4$ **unfolding** $\chi4-def\ ssv2\ ss2$ **by** *auto*

have $niss2$: $\neg isIntO\ ss2$
by (*metis list-all-append list-all-simps(1) ninter4 tr2l*)

have $rss2$: *reachO ss2* **and** $rssv12$: *reachV ssv1 reachV ssv2*
using $Opt.reach-validFromS-reach\ r\ ss2l\ tr2NE\ v4$ **apply** *blast*
unfolding $ssv1\ ssv2$ **using** $r(3,4)$ **using** $\chi4$ **unfolding** $\chi4-def$
using $Van.reach-validFromS-reach\ ssv1\ ssv2\ ssv1l\ ssv2l$ **by** *auto metis+*

have Δ : $\Delta\ ww\ ww1\ ww2\ s1\ ss2\ statA\ ssv1\ ssv2\ statOO$
using $\chi4$ **unfolding** $\chi4-def\ ss2[symmetric]\ ssv1[symmetric]\ ssv2[symmetric]$ **by** *auto*

have $s13'$: *validTransO (ss2, s2')*
by (*simp add: s24' ss2*)

note $vs24 = s24'$ [*unfolded ss2[symmetric]*]
obtain $w1'\ w2'\ trv1'\ trv2'\ statO'$ **where**
 $\chi4a$: $\chi4a\ \Delta\ \infty\ ww1\ ww2\ w1'\ w2'\ s1\ ss2\ s2'\ statA\ ssv1\ trv1'\ ssv2\ trv2'\ statO'$
using $unwindCond-ex-\chi4a-getSec$ [$OF\ unwind\ \Delta\ r(1)\ rss2\ rssv12\ s13'\ niss2\ iss2\ issv2\ gissv24$]
by *blast*

have $trv12'NE$: $trv1' \neq []\ trv2' \neq []$ **using** $\chi4a$ **unfolding** $\chi4a-def$ **by** *auto*

have [simp]: $\text{Van.O (butlast } trv1 \text{ @ } trv1') = \text{Van.O } trv1 \text{ @ Van.O } trv1'$
using $trv12'NE$ **unfolding** $\chi_4\text{-def}$ Van.O.map-filter Opt.O.map-filter **apply**(subst butlast-append) **by** simp

have [simp]: $\text{Van.O (butlast } trv2 \text{ @ } trv2') = \text{Van.O } trv2 \text{ @ Van.O } trv2'$
using $trv12'NE$ **unfolding** $\chi_4\text{-def}$ Van.O.map-filter Opt.O.map-filter **apply**(subst butlast-append) **by** simp

have $\text{Van.A } trv1' = \text{Van.A } trv2'$ **using** χ_4a **unfolding** $\chi_4a\text{-def}$ **by** auto
moreover **have** $\text{length (Van.O } trv1') = \text{length (Van.A } trv1') \wedge \text{length (Van.O } trv2') = \text{length (Van.A } trv2')$
unfolding Van.A.map-filter Van.O.map-filter **by** auto
ultimately **have** $\text{length (Van.O } trv1') = \text{length (Van.O } trv2')$ **by** auto
hence [simp]: $\text{Van.O } trv1 \text{ @ Van.O } trv1' = \text{Van.O } trv2 \text{ @ Van.O } trv2' \longleftrightarrow$
 $\text{Van.O } trv1 = \text{Van.O } trv2 \wedge \text{Van.O } trv1' = \text{Van.O } trv2'$ **by** auto

show ?thesis
apply(rule $exI[of - w1']$) **apply**(rule $exI[of - w2']$)
apply(rule $exI[of - butlast } trv1 \text{ @ } trv1']$) **apply**(rule $exI[of - butlast } trv2 \text{ @ } trv2']$)
apply(rule $exI[of - statO']$)
unfolding $\chi_4b\text{-def}$ **apply**(intro conjI)
subgoal **using** χ_4 χ_4a **unfolding** $\chi_4\text{-def}$ $\chi_4a\text{-def}$ **by** auto
subgoal **using** χ_4 χ_4a **unfolding** $\chi_4\text{-def}$ $\chi_4a\text{-def}$ **by** auto
subgoal **using** χ_4 χ_4a **unfolding** $\chi_4\text{-def}$ $\chi_4a\text{-def}$
by simp (metis Simple-Transition-System.fromS-eq-Nil Van.toS-Nil Van.toS-fromS-nonSingl

$\text{diff-add-inverse2 linorder-not-less order-le-less-trans trans-less-add2 zero-less-diff}$)

subgoal **using** χ_4 χ_4a **unfolding** $\chi_4\text{-def}$ $\chi_4a\text{-def}$ ssv1
using $\text{Van.validFromS-append}$ **by** auto
subgoal **using** χ_4 χ_4a **unfolding** $\chi_4\text{-def}$ $\chi_4a\text{-def}$ ssv2
using $\text{Van.validFromS-append}$ **by** auto
subgoal **using** χ_4 χ_4a **unfolding** $\chi_4\text{-def}$ $\chi_4a\text{-def}$
by (simp add: butlast-append)
subgoal **using** χ_4 χ_4a **unfolding** $\chi_4\text{-def}$ $\chi_4a\text{-def}$ **unfolding** Van.S.map-filter
 Opt.S.map-filter
apply(subst $tr2l$) **apply**(subst butlast-append)
by simp (metis Opt.S.map-filter Opt.S.eq-Nil-iff Van.S.map-filter Van.S.eq-Nil-iff
 $\text{sec}(1)$)
subgoal **using** χ_4 χ_4a **unfolding** $\chi_4\text{-def}$ $\chi_4a\text{-def}$ Van.A.map-filter Opt.A.map-filter

apply(subst $trv1l$) **apply**(subst $trv2l$)
apply(subst butlast-append) **apply** simp **apply**(subst butlast-append) **by** simp
subgoal **using** χ_4a $trv12'NE$ $tr2NE$ **unfolding** $\chi_4a\text{-def}$ lastt-def **by** simp .
qed

lemma unwindCond-ex- χ_4b -aux2:
assumes unwind: unwindCond Δ

and Δ : $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$ **and**
 r : $reachO s1 reachO s2 reachV sv1 reachV sv2$
and $v4'$: $Opt.validFromS s2 (tr2 @ [s2',s2'])$
and $nis1$: $\neg isIntO s1$ **and** $nis2$: $\neg isIntO s2$
and $ninter4'$: $never isIntO (tr2 @ [s2',s2'])$
and sec : $never isSecO tr2 isSecO s2'$
shows $\exists w1' w2' trv1 trv2 statOO. \chi4b \Delta \infty w1 w2 w1' w2' s1 s2 (tr2 @ [s2',s2'])$
 $statA sv1 trv1 sv2 trv2 statOO$
proof –
have 0 : $lastt s2 (tr2 \#\# s2') = s2'$
unfolding $lastt-def$ **by** $auto$
show $?thesis$
using $unwindCond-ex-\chi4b-aux[OF unwind \Delta r, of tr2 \#\# s2', unfolded 0, simplified]$
using $assms$ **by** $auto$
qed

definition $\chi4' \Delta w1 w2 w1' (w2'::enat) s1 s2 tr2 s2' s2'' statAA sv1 trv1 sv1''$
 $sv2 trv2 sv2'' statOO \equiv$
 $Van.validFromS sv1 (trv1 \#\# sv1'') \wedge Van.validFromS sv2 (trv2 \#\# sv2'') \wedge$
 $never isSecV (butlast (trv1 \#\# sv1'')) \wedge$
 $Van.S (trv2 \#\# sv2'') = Opt.S ((tr2 \#\# s2') \#\# s2'') \wedge$
 $Van.A (trv1 \#\# sv1'') = Van.A (trv2 \#\# sv2'') \wedge$
 $trv2 \neq [] \wedge (trv1 \neq [] \vee w1' < w1) \wedge$
 $\Delta \infty w1' w2' s1 s2'' statAA sv1'' sv2'' statOO$

proposition $unwindCond-ex-\chi4'$:
assumes $unwind$: $unwindCond \Delta$
and Δ : $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$ **and**
 r : $reachO s1 reachO s2 reachV sv1 reachV sv2$
and $v4'$: $Opt.validFromS s2 ((tr2 \#\# s2') \#\# s2'')$
and $nis1$: $\neg isIntO s1$ **and** $nis2$: $\neg isIntO s2$
and $ninter4'$: $never isIntO ((tr2 \#\# s2') \#\# s2'')$
and sec : $never isSecO tr2 isSecO s2'$
shows $\exists w1' w2' trv1 sv1'' trv2 sv2'' statOO. \chi4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2'$
 $s2'' statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$
using $unwindCond-ex-\chi4b-aux2[unfolded \varphi-def, unfolded lastt-snoc lastt-snoc2 append-snoc2, OF assms]$
unfolding $\chi4b-def$ **apply**($elim exE$) **subgoal for** $w1' w2' trv1 trv2 statOO$
apply($cases trv1$ $rule: rev-cases$)
subgoal by $auto$
subgoal for $trv1' sv1''$ **apply**($cases trv2$ $rule: rev-cases$)
subgoal by $auto$
subgoal for $trv2' sv2''$ **unfolding** $\chi4'-def$
apply($rule exI[of - w1']$) **apply**($rule exI[of - w2']$)
apply($rule exI[of - trv1']$) **apply**($rule exI[of - sv1'']$)

apply(rule exI[of - trv2']) **apply**(rule exI[of - sv2'])
apply(rule exI[of - statOO])
by simp (metis Opt.S.Nil-iff Opt.S.eq-Nil-iff(1) Van.S.simps(4) butlast-append
list.discI list-all-append sec(2) self-append-conv2) . . .

definition $\omega_4 \Delta w_1 w_2 w_1' (w_2'::\text{enat}) s_1 s_2 s_2' \text{statAA} sv_1 trv_1 sv_1' sv_2 trv_2 sv_2' \text{statOO} \equiv$
Van.validFromS sv1 (trv1 ## sv1') \wedge Van.validFromS sv2 (trv2 ## sv2') \wedge
never isSecV trv1 \wedge never isSecV trv2 \wedge
Van.A (trv1 ## sv1') = Van.A (trv2 ## sv2') \wedge
(trv1 \neq [] \vee w1' < w1) \wedge (trv2 \neq [] \vee w2' < w2) \wedge
 $\Delta \infty w_1' w_2' s_1 s_2' \text{statAA} sv_1' sv_2' \text{statOO}$

proposition unwindCond-ex- ω_4 :
assumes unwind: unwindCond Δ
and Δ : $\Delta w w_1 w_2 s_1 s_2 \text{statA} sv_1 sv_2 \text{statO}$
and r34: reachO s1 reachO s2 **and** r12: reachV sv1 reachV sv2
and nis1: \neg isIntO s1
and v4: validTransO (s2,s2')
and nis2: \neg isIntO s2 \neg isIntO s2' \neg isSecO s2
shows $\exists w_1' w_2' trv_1 sv_1' trv_2 sv_2' \text{statOO}. \omega_4 \Delta w_1 w_2 w_1' w_2' s_1 s_2 s_2' \text{statA} sv_1 trv_1 sv_1' sv_2 trv_2 sv_2' \text{statOO}$
using Δ r12
proof(induction w arbitrary: w1 w2 sv1 sv2 statO rule: less-induct)
case (less w w1 w2 sv1 sv2 statO)
note $\Delta = \langle \Delta w w_1 w_2 s_1 s_2 \text{statA} sv_1 sv_2 \text{statO} \rangle$
note r12 = less.prem(2,3)
note r1 = r12(1) **note** r2 = r12(2)
note r = r34 r12

have f34: finalO s1 = finalO s2 \wedge finalV sv1 = finalO s1 \wedge finalV sv2 = finalO s2
using Δ unwind[unfolded unwindCond-def] r **by** auto

have proact-match: ($\exists v < w. \text{proact } \Delta v w_1 w_2 s_1 s_2 \text{statA} sv_1 sv_2 \text{statO}$) \vee match
 $\Delta w_1 w_2 s_1 s_2 \text{statA} sv_1 sv_2 \text{statO}$
using Δ unwind[unfolded unwindCond-def] r **by** auto
show ?case **using** proact-match **proof** safe
fix v **assume** v: v < w
assume proact $\Delta v w_1 w_2 s_1 s_2 \text{statA} sv_1 sv_2 \text{statO}$
thus ?thesis **unfolding** proact-def **proof** safe
assume sv1: \neg isSecV sv1 \neg isIntV sv1 **and** move-1 $\Delta v w_1 w_2 s_1 s_2 \text{statA} sv_1 sv_2 \text{statO}$
then obtain sv1' **where** 0: validTransV (sv1, sv1') **and** Δ : $\Delta v w_1 w_2 s_1 s_2 \text{statA} sv_1' sv_2 \text{statO}$
unfolding move-1-def **by** auto

```

have r1': reachV sv1' using r1 0 by (metis Van.reach.Step fst-conv snd-conv)
obtain w1' w2' trv1 sv1'' trv2 sv2' statOO where
ω4: ω4 Δ w1 w2 w1' w2' s1 s2 s2' statA sv1' trv1 sv1'' sv2 trv2 sv2' statOO

using less(1)[OF v Δ r1' r2] by auto
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
exI[of - sv1 # trv1]) apply(rule exI[of - sv1''])
apply(rule exI[of - trv2]) apply(rule exI[of - sv2'])
using ω4 0 sv1 unfolding ω4-def by auto
next
assume sv2: ¬ isSecV sv2 ¬ isIntV sv2 and move-2 Δ v w1 w2 s1 s2 statA
sv1 sv2 statO
then obtain sv2'
where 0: validTransV (sv2,sv2')
and Δ: Δ v w1 w2 s1 s2 statA sv1 sv2' statO
unfolding move-2-def by auto
have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
obtain w1' w2' trv1 sv1' trv2 sv2'' statOO where
ω4: ω4 Δ w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2' trv2 sv2'' statOO

using less(1)[OF v Δ r1 r2'] by auto
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - trv1]) apply(rule exI[of - sv1'])
apply(rule exI[of - sv2 # trv2]) apply(rule exI[of - sv2''])
using ω4 0 sv2 unfolding ω4-def by auto
next
assume sv1: ¬ isSecV sv1 and sv2: ¬ isSecV sv2 and
move-12 Δ v w1 w2 s1 s2 statA sv1 sv2 statO and
sv12: Van.eqAct sv1 sv2
then obtain sv1' sv2' statO'
where statO': statO' = sstatO' statO sv1 sv2
and 0: validTransV (sv1,sv1') validTransV (sv2,sv2')
and Δ: Δ v w1 w2 s1 s2 statA sv1' sv2' statO'
unfolding move-12-def by auto
have r1': reachV sv1' and r2': reachV sv2' using r1 r2 0
by (metis Van.reach.Step fst-conv snd-conv)+
obtain w1' w2' trv1 sv1'' trv2 sv2'' statOO where
ω4: ω4 Δ w1 w2 w1' w2' s1 s2 s2' statA sv1' trv1 sv1'' sv2' trv2 sv2'' statOO

using less(1)[OF v Δ r1' r2'] by auto
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv1''])
apply(rule exI[of - sv2 # trv2]) apply(rule exI[of - sv2''])
using ω4 0 sv1 sv2 sv12 unfolding ω4-def statO' by (auto simp: Van.eqAct-def)
qed
next
assume m: match Δ w1 w2 s1 s2 statA sv1 sv2 statO
have m: match2 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
match-def by auto

```


have $(\exists w1' w2'. w1' < w1 \wedge w2' < w2 \wedge \neg isSecO\ s2 \wedge \Delta \infty\ w1' w2' s1\ s2' statA\ sv1\ sv2\ statO) \vee$
 $(\exists w1' < w1. eqSec\ sv2\ s2 \wedge \neg isIntV\ sv2 \wedge match2-1\ \Delta\ w1' \infty\ s1\ s2\ s2' statA\ sv1\ sv2\ statO) \vee$
 $\neg isSecV\ sv1 \wedge eqSec\ sv2\ s2 \wedge Van.eqAct\ sv1\ sv2 \wedge match2-12\ \Delta\ \infty \infty s1\ s2\ s2' statA\ sv1\ sv2\ statO$
using $m\ v4\ nis2$ **unfolding** $match2-def$ **by** $auto$

thus $?thesis$

apply $(elim\ disjE\ exE)$

subgoal for $w1' w2'$ **apply** $(rule\ exI[of - w1'])$ **apply** $(rule\ exI[of - w2'])$

apply $(rule\ exI[of - []])$ **apply** $(rule\ exI[of - sv1])$

apply $(rule\ exI[of - []])$ **apply** $(rule\ exI[of - sv2])$

apply $(rule\ exI[of - statO])$ **unfolding** $\omega4-def$

by $auto$

subgoal for $w1'$ **apply** $(rule\ exI[of - w1'])$ **apply** $(rule\ exI[of - \infty])$

unfolding $match2-1-def$ **apply** $(elim\ conjE\ exE)$ **subgoal for** $sv2'$

apply $(rule\ exI[of - []])$ **apply** $(rule\ exI[of - sv1])$

apply $(rule\ exI[of - [sv2]])$ **apply** $(rule\ exI[of - sv2'])$

apply $(rule\ exI[of - statO])$

unfolding $\omega4-def$ **using** $nis2(3)$ **by** $(auto\ simp: eqSec-def)$.

subgoal **apply** $(rule\ exI[of - \infty])$ **apply** $(rule\ exI[of - \infty])$

unfolding $match2-12-def$ **apply** $(elim\ conjE\ exE)$ **subgoal for** $sv1' sv2'$

apply $(rule\ exI[of - [sv1]])$ **apply** $(rule\ exI[of - sv1'])$

apply $(rule\ exI[of - [sv2]])$ **apply** $(rule\ exI[of - sv2'])$

apply $(rule\ exI[of - sstatO' statO sv1 sv2])$

unfolding $\omega4-def$ **using** $nis2(3)$ **apply** $(auto\ simp: eqSec-def$

$sstatA'-def\ sstatO'-def\ lastt-def\ Van.eqAct-def) \dots$

qed

qed

definition $\varphi\varphi\ s1\ ltr1\ s2\ ltr2\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2' \equiv$

$ltr1 = lappend\ (l\ list\ of\ (tr1\ \#\#\ s1'))\ (s1''\ \$\ ltr1') \wedge$

$ltr2 = lappend\ (l\ list\ of\ (tr2\ \#\#\ s2'))\ (s2''\ \$\ ltr2') \wedge$

$Opt.validFromS\ s1\ ((tr1\ \#\#\ s1')\ \#\#\ s1'') \wedge Opt.validFromS\ s2\ ((tr2\ \#\#\ s2')\ \#\#\ s2'') \wedge$

$never\ isIntO\ tr1 \wedge never\ isIntO\ tr2 \wedge$

$isIntO\ s1' \wedge isIntO\ s2' \wedge getActO\ s1' = getActO\ s2' \wedge$

$Opt.lvalidFromS\ s1''\ (s1''\ \$\ ltr1') \wedge Opt.lcompletedFrom\ s1''\ (s1''\ \$\ ltr1') \wedge$

$Opt.lvalidFromS\ s2''\ (s2''\ \$\ ltr2') \wedge Opt.lcompletedFrom\ s2''\ (s2''\ \$\ ltr2') \wedge$

$Opt.lA\ (s1''\ \$\ ltr1') = Opt.lA\ (s2''\ \$\ ltr2')$

lemma *isIntO- $\varphi\varphi$* :
assumes *vltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1*
and *vltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2*
and *A: Opt.lA ltr1 = Opt.lA ltr2* **and** *inter3: \neg lnever isIntO ltr1*
shows $\exists tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'. \varphi\varphi s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'$
proof –
have *03: $\exists s \in \text{lset } ltr1. \text{isIntO } s$* **using** *inter3* **unfolding** *llist.pred-set* **by** *auto*
define *ttr1* **where** *ttr1: ttr1 \equiv ltakeUntil isIntO ltr1*
define *lltr1'* **where** *lltr1': lltr1' \equiv ldropUntil isIntO ltr1*
have *ltr1: ltr1 = lappend (llist-of ttr1) lltr1'*
unfolding *ttr1 lltr1' lappend-ltakeUntil-ldropUntil[OF 03]* **..**
have *13: ttr1 \neq [] \wedge never isIntO (butlast ttr1) \wedge isIntO (last ttr1)*
unfolding *ttr1*
using *ltakeUntil-last[OF 03] ltakeUntil-not-Nil[OF 03] ltakeUntil-never-butlast[OF 03]* **by** *simp*
then obtain *tr1 s1'* **where** *ttr1-eq: ttr1 = tr1 ## s1'*
using *rev-exhaust* **by** *blast*
hence *tr1s1': never isIntO tr1 isIntO s1'* **using** *13* **by** *auto*
have *lfinite ltr1 \implies s1' \neq llast ltr1*
by (*metis Opt.final-not-isInt Opt.lcompletedFrom-def llast-last-llist-of tr1s1'(2)*)
vltr1(2)
hence *ne: lltr1' \neq []*
using *ltr1* **unfolding** *ttr1-eq* **by** *auto*
then obtain *s1'' ltr1'* **where** *lltr1': lltr1' = s1'' \$ ltr1'*
by (*meson llist.exhaust*)
have [*simp*]: *filter isIntO tr1 = []*
by (*metis never-Nil-filter tr1s1'(1)*)
have *clltr1': Opt.lcompletedFrom s1 lltr1'*
by (*metis Opt.lcompletedFrom-def lfinite-lappend lfinite-llist-of llast-lappend-LCons*
llast-last-llist-of lltr1' ltr1 ne vltr1(2))

have *inter4: \neg lnever isIntO ltr2* **using** *A inter3*
by (*metis Opt.lA.eq-LNil-iff Opt.lO Opt.lO.eq-LNil-iff lfiltermap-LNil-never*
lfiltermap-lmap-lfilter vltr1(2) vltr2(2))
have *04: $\exists s \in \text{lset } ltr2. \text{isIntO } s$* **using** *inter4* **unfolding** *llist.pred-set* **by** *auto*
define *ttr2* **where** *ttr2: ttr2 \equiv ltakeUntil isIntO ltr2*
define *lltr2'* **where** *lltr2': lltr2' \equiv ldropUntil isIntO ltr2*
have *ltr2: ltr2 = lappend (llist-of ttr2) lltr2'*
unfolding *ttr2 lltr2' lappend-ltakeUntil-ldropUntil[OF 04]* **..**
have *14: ttr2 \neq [] \wedge never isIntO (butlast ttr2) \wedge isIntO (last ttr2)*
unfolding *ttr2*
using *ltakeUntil-last[OF 04] ltakeUntil-not-Nil[OF 04] ltakeUntil-never-butlast[OF 04]* **by** *simp*
then obtain *tr2 s2'* **where** *ttr2-eq: ttr2 = tr2 ## s2'*
using *rev-exhaust* **by** *blast*
hence *tr2s2': never isIntO tr2 isIntO s2'* **using** *14* **by** *auto*

have $lfinite\ ltr2 \implies s2' \neq llast\ ltr2$
by (*metis* *Opt.final-not-isInt* *Opt.lcompletedFrom-def* *llast-last-llist-of* *tr2s2'(2)* *vltr2(2)*)
hence $ne: lltr2' \neq []$
using *ltr2* **unfolding** *ttr2-eq* **by** *auto*
then obtain $s2''\ ltr2'$ **where** $lltr2': lltr2' = s2''\ \$\ ltr2'$
by (*meson* *llist.exhaust*)
have [*simp*]: $filter\ isIntO\ tr2 = []$
by (*metis* *never-Nil-filter* *tr2s2'(1)*)
have $clltr2': Opt.lcompletedFrom\ s2\ lltr2'$
by (*metis* *Opt.lcompletedFrom-def* *lfinite-lappend* *lfinite-llist-of* *llast-lappend-LCons* *llast-last-llist-of* *lltr2' ltr2 ne vltr2(2)*)

have $AA: Opt.lA\ lltr1' = Opt.lA\ lltr2'$
unfolding $Opt.lA[OF\ clltr1']\ Opt.lA[OF\ clltr2']$
using $A[unfolding\ Opt.lA[OF\ vltr1(2)]\ Opt.lA[OF\ vltr2(2)]]\ tr1s1'\ tr2s2'$
unfolding *ltr1* *ltr2* *ttr1-eq* *ttr2-eq*
unfolding *lfilter-lappend-llist-of* **by** *simp*

show *?thesis* **apply**(*rule* $exI[of - tr1]$) **apply**(*rule* $exI[of - s1']$)
apply(*rule* $exI[of - s1']$) **apply**(*rule* $exI[of - ltr1']$)
apply(*rule* $exI[of - tr2]$) **apply**(*rule* $exI[of - s2']$)
apply(*rule* $exI[of - s2']$) **apply**(*rule* $exI[of - ltr2']$)
unfolding $\varphi\varphi\text{-def}$ **apply**(*intro* *conjI*)
subgoal **unfolding** *ltr1* *ttr1-eq* *lltr1' ..*
subgoal **unfolding** *ltr2* *ttr2-eq* *lltr2' ..*
subgoal **using** *vltr1(1)* **unfolding** *ltr1* *ttr1-eq* *lltr1'*
by (*simp* *add: Opt.lvalidFromS-lappend-finite* *lappend-llist-of-LCons*)
subgoal **using** *vltr2(1)* **unfolding** *ltr2* *ttr2-eq* *lltr2'*
by (*simp* *add: Opt.lvalidFromS-lappend-finite* *lappend-llist-of-LCons*)
subgoal **using** *tr1s1'* **by** *simp*
subgoal **using** *tr2s2'* **by** *simp*
subgoal **using** *tr1s1'* **by** *simp*
subgoal **using** *tr2s2'* **by** *simp*
subgoal **using** $A[unfolding\ Opt.lA[OF\ vltr1(2)]\ Opt.lA[OF\ vltr2(2)]]\ tr1s1'\ tr2s2'$
unfolding *ltr1* *ttr1-eq* *ltr2* *ttr2-eq* *lltr1' lltr2'*
unfolding *lfilter-lappend-llist-of* **by** *simp*
subgoal **using** *vltr1(1)* **unfolding** *ltr1* *ttr1-eq* *lltr1'*
using *Opt.lvalidFromS-lappend-LCons* **by** *blast*
subgoal **using** *vltr1(2)* **unfolding** *ltr1* *ttr1-eq* *lltr1'*
by (*metis* *Opt.lcompletedFrom-def* *lfinite-lappend* *lfinite-llist-of* *llast-lappend-LCons* *llast-last-llist-of* *llist.distinct(1)*)
subgoal **using** *vltr2(1)* **unfolding** *ltr2* *ttr2-eq* *lltr2'*
using *Opt.lvalidFromS-lappend-LCons* **by** *blast*
subgoal **using** *vltr2(2)* **unfolding** *ltr2* *ttr2-eq* *lltr2'*
by (*metis* *Opt.lcompletedFrom-def* *lfinite-lappend* *lfinite-llist-of* *llast-lappend-LCons* *llast-last-llist-of* *llist.distinct(1)*)

subgoal using *AA unfolding* $lltr1' lltr2' . .$
qed

definition $\chi\chi$ $s1$ $ltr1$ $tr1$ $s1'$ $s1''$ $ltr1' \equiv$
 $ltr1 = lappend$ (l list-of ($tr1 \## s1'$)) ($s1'' \$ ltr1'$) \wedge
 $Opt.validFromS$ $s1$ ($(tr1 \## s1') \## s1''$) \wedge
 $never$ $isIntO$ $tr1$ $\wedge \neg isIntO$ $s1'$ $\wedge \neg isIntO$ $s1''$ \wedge
 $never$ $isSecO$ $tr1$ $\wedge isSecO$ $s1'$ \wedge
 $Opt.lvalidFromS$ $s1''$ ($s1'' \$ ltr1'$) $\wedge Opt.lcompletedFrom$ $s1''$ ($s1'' \$ ltr1'$)

lemma $isSecO$ - $\chi\chi$:

assumes $vltr1$: $Opt.lvalidFromS$ $s1$ $ltr1$ $Opt.lcompletedFrom$ $s1$ $ltr1$

and $inter$: $lnever$ $isIntO$ $ltr1$ **and** $isec$: $\neg lnever$ $isSecO$ $ltr1$

shows $\exists tr1$ $s1'$ $s1''$ $ltr1'$. $\chi\chi$ $s1$ $ltr1$ $tr1$ $s1'$ $s1''$ $ltr1'$

proof –

have 0 : $\exists s \in lset$ $ltr1$. $isSecO$ s **using** $isec$ **unfolding** l list.pred-set by *auto*
define $ttr1$ **where** $ttr1$: $ttr1 \equiv ltakeUntil$ $isSecO$ $ltr1$
define $lltr1'$ **where** $lltr1'$: $lltr1' \equiv ldropUntil$ $isSecO$ $ltr1$
have $ltr1$: $ltr1 = lappend$ (l list-of $ttr1$) $lltr1'$
unfolding $ttr1$ $lltr1'$ $lappend$ - $ltakeUntil$ - $ldropUntil$ [*OF* 0] ..
have 1 : $ttr1 \neq []$ $\wedge never$ $isSecO$ ($butlast$ $ttr1$) $\wedge isSecO$ ($last$ $ttr1$)
unfolding $ttr1$
using $ltakeUntil$ - $last$ [*OF* 0] $ltakeUntil$ - not - Nil [*OF* 0] $ltakeUntil$ - $never$ - $butlast$ [*OF* 0] by *simp*
then obtain $tr1$ $s1'$ **where** $ttr1$ - eq : $ttr1 = tr1 \## s1'$
using rev - $exhaust$ by *blast*
hence $tr1s1'$: $never$ $isSecO$ $tr1$ $isSecO$ $s1'$ **using** 1 by *auto*
have 2 : $never$ $isIntO$ $tr1$ $\wedge \neg isIntO$ $s1'$ $\wedge lnever$ $isIntO$ $lltr1'$
using $inter$ **unfolding** $ltr1$ $ttr1$ - eq
unfolding l list-all- $lappend$ - l list-of l ist-all- $append$ by *simp*
have $lfinite$ $ltr1 \implies s1' \neq llast$ $ltr1$
by ($metis$ $Opt.final$ - not - $isSec$ $Opt.lcompletedFrom$ - def $llast$ - $last$ - l list-of $tr1s1'$)(2)
 $vltr1$ (2)
hence ne : $lltr1' \neq []$
using $ltr1$ **unfolding** $ttr1$ - eq by *auto*
then obtain $s1''$ $ltr1'$ **where** $lltr1'$: $lltr1' = s1'' \$ ltr1'$
by ($meson$ l list. $exhaust$)
show $?thesis$ **apply**($rule$ exI [of - $tr1$]) **apply**($rule$ exI [of - $s1'$])
apply($rule$ exI [of - $s1''$]) **apply**($rule$ exI [of - $ltr1'$])
unfolding $\chi\chi$ - def **apply**($intro$ $conjI$)
subgoal unfolding $ltr1$ $ttr1$ - eq $lltr1'$..
subgoal using $vltr1$ (1) **unfolding** $ltr1$ $ttr1$ - eq $lltr1'$
by ($simp$ add : $Opt.lvalidFromS$ - $lappend$ - $finite$ $lappend$ - l list-of- $LCons$)
subgoal using 2 by *simp*
subgoal using 2 by *simp*
subgoal using 2 **unfolding** $lltr1'$ by *simp*
subgoal using $tr1s1'$ by *simp*
subgoal using $tr1s1'$ by *simp*
subgoal using $vltr1$ (1) **unfolding** $ltr1$ $ttr1$ - eq $lltr1'$

```

using Opt.lvalidFromS-lappend-LCons by blast
subgoal using vltr1(2) unfolding ltr1 ttr1-eq lltr1'
by (metis Opt.lcompletedFrom-def ne lfinite-lappend
     lfinite-llist-of llast-lappend-LCons llast-last-llist-of lltr1') .
qed

```

```

type-synonym ('stA,'stO) tuple34 =
  enat × enat ×
  'stA × 'stA llist ×
  'stA × 'stA llist ×
  status ×
  'stO × 'stO × status

```

```

type-synonym ('stA,'stO) tuple12 =
  'stO list × 'stO × 'stO list × 'stO × status × status

```

context

```

fixes  $\Delta$  :: enat ⇒ enat ⇒ enat ⇒ 'stateO ⇒ 'stateO ⇒ status ⇒ 'stateV ⇒
'stateV ⇒ status ⇒ bool
begin

```

```

fun isn :: turn × ('stateO,'stateV) tuple34 ⇒ bool

```

where

```

isn (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ←→ ltr1 = [] ∧ ltr2 = []

```

fun *h-t* ::

```

turn × ('stateO,'stateV)tuple34 ⇒
  ('stateO,'stateV)tuple12 ×
  turn × ('stateO,'stateV)tuple34

```

where

```

h-t (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) =
  (if trn = L
   then if lnever isSecO ltr1
    then let (s1',ltr1') = (lhd (ltl ltr1), ltl ltr1)
     in let (w1',w2',trv1,sv1',trv2,sv2',statOO) =
        (SOME k. case k of (w1',w2',trv1,sv1',trv2,sv2',statOO) ⇒
           ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2' statOO)
     in ((trv1,sv1',trv2,sv2',statA,statOO),
        (if trv1 = [] then L else R,
         w1',w2',s1',ltr1',s2,ltr2,statA,sv1',sv2',statOO))
   else
    let (tr1,s1',s1'',ltr1') =
        (SOME k. case k of (tr1,s1',s1'',ltr1') ⇒

```

```

     $\chi\chi$  s1 ltr1 tr1 s1' s1'' ltr1')
  in let (w1',w2',trv1,sv1'',trv2,sv2'',statOO) =
    (SOME k'. case k' of (w1',w2',trv1,sv1'',trv2,sv2'',statOO)  $\Rightarrow$ 
       $\chi\chi^3$ '  $\Delta$  w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2
    sv2'' statOO)
  in ((trv1,sv1'',trv2,sv2'',statA,statOO),
    (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))
  —
  else if lnever isSecO ltr2
  then let (s2',ltr2') = (lhd (ltl ltr2), ltl ltr2)
  in let (w1',w2',trv1,sv1',trv2,sv2',statOO) =
    (SOME k. case k of (w1',w2',trv1,sv1',trv2,sv2',statOO)  $\Rightarrow$ 
       $\omega_4$   $\Delta$  w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2' statOO)
  in ((trv1,sv1',trv2,sv2',statA,statOO),
    (if trv2 = [] then R else L,
      w1',w2',s1,ltr1,s2',ltr2',statA,sv1',sv2',statOO))
  else
  let (tr2,s2',s2'',ltr2') =
    (SOME k. case k of (tr2,s2',s2'',ltr2')  $\Rightarrow$ 
       $\chi\chi$  s2 ltr2 tr2 s2' s2'' ltr2')
  in let (w1',w2',trv1,sv1'',trv2,sv2'',statOO) =
    (SOME k'. case k' of (w1',w2',trv1,sv1'',trv2,sv2'',statOO)  $\Rightarrow$ 
       $\chi\chi^4$ '  $\Delta$  w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2 trv2
    sv2'' statOO)
  in ((trv1,sv1'',trv2,sv2'',statA,statOO),
    (L,w1',w2',s1, ltr1, s2'',s2'' $ ltr2',statA,sv1'',sv2'',statOO))
)

```

declare *h-t.simps*[simp del]

definition *h* \equiv *fst* o *h-t*

definition *t* \equiv *snd* o *h-t*

fun *econd* **where** *econd* (*trn*,*w1*,*w2*,*s1*,*ltr1*,*s2*,*ltr2*,*statA*,*sv1*,*sv2*,*statO*) =
 (*llength* *ltr1* \leq *Suc* 0 \vee *llength* *ltr2* \leq *Suc* 0)

fun *e* **where** *e* (*trn*,*w1*,*w2*,*s1*,*ltr1*,*s2*,*ltr2*,*statA*,*sv1*,*sv2*,*statO*) = [[([*sv1*],*sv1*],[*sv2*],*sv2*,*statA*,*statO*)]

definition *f* :: *turn* \times ('*stateO*','*stateV*)*tuple34* \Rightarrow ('*stateO*','*stateV*)*tuple12* *llist*

where *f* \equiv *ccorec-llist* *isn* *h* *econd* *e* *t*

lemma *f-LNil*:

ltr1 = [] \Rightarrow *ltr2* = [] \Rightarrow *f* (*trn*,*w1*,*w2*,*s1*,*ltr1*,*s2*,*ltr2*,*statA*,*sv1*,*sv2*,*statO*) = []

unfolding *f-def* **apply**(*subst* *llist-ccorec*(1)) **by** *auto*

lemma *f-length-1*:

assumes $ltr1 \neq [] \vee ltr2 \neq [] \wedge llength\ ltr1 \leq Suc\ 0 \vee llength\ ltr2 \leq Suc\ 0$
shows $f\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = [([sv1], sv1, [sv2], sv2, statA, statO)]$

using *assms* **unfolding** *f-def* **apply**(*subst llist-ccorec(2)*)
subgoal unfolding *e.simps lnull-def* **by** *auto*
subgoal by *auto*
subgoal unfolding *econd.simps* **by** *simp* .

lemma *f-length-ge1*:

assumes $llength\ ltr1 > Suc\ 0 \wedge llength\ ltr2 > Suc\ 0$

shows $f\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$

$LCons\ (h\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO))\ (f\ (t\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, sta$
 $))$

proof –

show *?thesis* **using** *assms* **unfolding** *f-def* **apply**(*subst llist-ccorec(2)*)

subgoal unfolding *e.simps lnull-def* **by** *auto*

subgoal by *auto*

subgoal unfolding *econd.simps* **by** *auto* .

qed

definition *lltrv1* :: $turn \times ('stateO, 'stateV)tuple34 \Rightarrow 'stateV\ llist$ **where**

$lltrv1\ trn-tp = lconcat\ (lmap\ (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO). llist-of\ trv1)$
 $(f\ trn-tp))$

definition *then1* :: $turn \times ('stateO, 'stateV)tuple34 \Rightarrow nat$ **where**

$then1\ trn-tp = firstNC\ (lmap\ (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO). trv1)\ (f\ trn-tp))$

definition *lltrv2* :: $turn \times ('stateO, 'stateV)tuple34 \Rightarrow 'stateV\ llist$ **where**

$lltrv2\ trn-tp = lconcat\ (lmap\ (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO). llist-of\ trv2)$
 $(f\ trn-tp))$

definition *then2* :: $turn \times ('stateO, 'stateV)tuple34 \Rightarrow nat$ **where**

$then2\ trn-tp = firstNC\ (lmap\ (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO). trv2)\ (f\ trn-tp))$

lemma *lltrv1-ne-imp*:

assumes $lltrv1\ trn-tp \neq []$

shows $\exists\ trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO. (trv1, sv1'', trv2, sv2'', statAA, statOO)$
 $\in\ lset\ (f\ trn-tp) \wedge$

$trv1 \neq []$

using *assms* **unfolding** *lltrv1-def* **unfolding** *lconcat-eq-LNil-iff* **by** *force*

lemma *lltrv2-ne-imp*:

assumes $lltrv2\ trn-tp \neq []$

shows $\exists\ trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO. (trv1, sv1'', trv2, sv2'', statAA, statOO)$
 $\in\ lset\ (f\ trn-tp) \wedge$

$trv2 \neq []$

using *assms* unfolding *lltrv2-def* unfolding *lconcat-eq-LNil-iff* by *force*

lemma *lltrv1-LNil[simp]*:

$ltr1 = [] \implies ltr2 = [] \implies lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)$
 $= []$

unfolding *lltrv1-def f-LNil* by *simp*

lemma *lltrv2-LNil[simp]*:

$ltr1 = [] \implies ltr2 = [] \implies lltrv2 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)$
 $= []$

unfolding *lltrv2-def f-LNil* by *simp*

lemma *lltrv1-lnever[simp]*:

assumes $ltr1 \neq [] \vee ltr2 \neq []$ $llength\ ltr1 \leq Suc\ 0 \vee llength\ ltr2 \leq Suc\ 0$

shows $lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = [[sv1]]$

unfolding *lltrv1-def* using *f-length-1[OF assms]* by *auto*

lemma *lltrv2-lnever[simp]*:

assumes $ltr1 \neq [] \vee ltr2 \neq []$ $llength\ ltr1 \leq Suc\ 0 \vee llength\ ltr2 \leq Suc\ 0$

shows $lltrv2 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = [[sv2]]$

unfolding *lltrv2-def* using *f-length-1[OF assms]* by *auto*

lemma *h-t-lnever-L*:

assumes *unw*: *unwindCond* Δ

and Δ : $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*

and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1*

and *l'*: *lnever isIntO* *ltr1* \neg *isIntO* *s2*

and *len*: $llength\ ltr1 > Suc\ 0$ $llength\ ltr2 > Suc\ 0$

and *l*: *trn* = *L* *lnever isSecO* *ltr1*

shows $\exists w1'\ w2'\ s1'\ ltr1'\ trv1\ sv1'\ trv2\ sv2'\ statOO.$

$ltr1 = s1 \$ ltr1' \wedge validTransO (s1, s1') \wedge$

$Opt.lvalidFromS\ s1'\ ltr1' \wedge Opt.lcompletedFrom\ s1'\ ltr1' \wedge lnever\ isIntO\ ltr1' \wedge$

$\omega^3\ \Delta\ w1\ w2\ w1'\ w2'\ s1\ s1'\ s2\ statA\ sv1\ trv1\ sv1'\ sv2\ trv2\ sv2'\ statOO \wedge$

$h-t (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$

$((trv1, sv1', trv2, sv2', statA, statOO),$

$(if\ trv1 = []\ then\ L\ else\ R,$

$w1', w2', s1', ltr1', s2, ltr2, statA, sv1', sv2', statOO))$

proof –

have *s1*: $\neg isIntO\ s1$ using *l' ltr1*

by (*metis* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-def* *lfinite-LNil* *lhd-LCons* *llist.exhaust* *llist.pred-inject*(2))

obtain *ltr1'* **where** *ltr13*: *ltr1* = *s1* \$ *ltr1'*

by (*metis* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-def* *lfinite-LNil* *llist.exhaust-sel* *ltr1*(1) *ltr1*(2))

hence *ltr1'*: *ltr1'* = *ltl* *ltr1* **by** *auto*

have *ltr1'ne*: *ltr1'ne* ≠ [] **using** *len*(1) **unfolding** *ltr13*

by (*metis* *One-nat-def* *llength-LCons* *llength-LNil* *one-eSuc* *one-enat-def* *order-less-irrefl*)

define *s1'* **where** *s1'*: *s1'* = *lhd* (*ltl* *ltr1*)

have *v3*: *validTransO* (*s1*, *s1'*) **and** *vv3*: *Opt.lvalidFromS* *s1'* *ltr1'* *Opt.lcompletedFrom* *s1'* *ltr1'*

using *ltr1* *ltr1'ne* **unfolding** *ltr13* *s1'*

by (*metis* *Opt.lcompletedFrom-LCons* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-Cons-iff* *ltr1'* *ltr13*)+

have *is1'*: ¬ *isIntO* *s1'* **and** *lnever* *isIntO* *ltr1'*

using *l'(1)* **unfolding** *ltr13*

by (*metis* *llist.exhaust-sel* *llist.pred-inject*(2) *ltr1'* *ltr1'ne* *s1'*)+

have *iss1*: ¬ *isSecO* *s1*

using *l(2)* *ltr13* **by** *auto*

obtain *w1'* *w2'* *trv1* *sv1'* *trv2* *sv2'* *statOO*

where *ω3*: *ω3* Δ *w1* *w2* *w1'* *w2'* *s1* (*lhd* (*ltl* *ltr1*)) *s2* *statA* *sv1* *trv1* *sv1'* *sv2* *trv2* *sv2'* *statOO*

using *unwindCond-ex-ω3*[*OF* *unw* Δ *r* *v3* *s1* *is1'* *iss1* *l'(2)*] *s1'* **by** *auto*

define *tp'* **where**

tp' = (*SOME* *k'*. *case* *k'* of (*w1'*, *w2'*, *trv1*, *sv1'*, *trv2*, *sv2'*, *statOO*) ⇒

ω3 Δ *w1* *w2* *w1'* *w2'* *s1* (*lhd* (*ltl* *ltr1*)) *s2* *statA* *sv1* *trv1* *sv1'* *sv2* *trv2* *sv2'* *statOO*)

have *1*: *case* *tp'* of (*w1'*, *w2'*, *trv1*, *sv1'*, *trv2*, *sv2'*, *statOO*) ⇒

ω3 Δ *w1* *w2* *w1'* *w2'* *s1* *s1'* *s2* *statA* *sv1* *trv1* *sv1'* *sv2* *trv2* *sv2'* *statOO*

using *ω3* **unfolding** *tp'-def* *s1'* **apply**– **apply**(*rule* *someI-ex*)

apply(*rule* *exI*[of - (*w1'*, *w2'*, *trv1*, *sv1'*, *trv2*, *sv2'*, *statOO*)]) **by** *auto*

obtain *w1'* *w2'* *trv1* *sv1'* *trv2* *sv2'* *statOO* **where**

tp': *tp'* = (*w1'*, *w2'*, *trv1*, *sv1'*, *trv2*, *sv2'*, *statOO*) **by**(*cases* *tp'*, *auto*)

have *ω3*: *ω3* Δ *w1* *w2* *w1'* *w2'* *s1* *s1'* *s2* *statA* *sv1* *trv1* *sv1'* *sv2* *trv2* *sv2'* *statOO*

using *1* **unfolding** *tp'* **by** *auto*

show *?thesis*

```

apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1']) apply(rule exI[of - ltr1'])
apply(rule exI[of - trv1]) apply(rule exI[of - sv1'])
apply(rule exI[of - trv2]) apply(rule exI[of - sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal using len l unfolding h-t.simps apply simp
  unfolding tp'-def[symmetric] tp' s1' ltr1' by simp .
qed

```

```

lemma lltrv1-lltrv2-lnever-L:
assumes unw: unwindCond Δ
and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and l': lnever isIntO ltr1  $\neg$  isIntO s2
and len: llength ltr1 > Suc 0 llength ltr2 > Suc 0
and l: trn = L lnever isSecO ltr1
shows  $\exists w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO.$ 
   $ltr1 = s1 \$ ltr1' \wedge validTransO (s1, s1') \wedge$ 
   $Opt.lvalidFromS s1' ltr1' \wedge Opt.lcompletedFrom s1' ltr1' \wedge lnever isIntO ltr1' \wedge$ 

   $\omega^3 \Delta w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2' statOO \wedge$ 
   $lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$ 
   $lappend (l\list-of trv1) (lltrv1 (if trv1 = [] then L else R,$ 
   $w1', w2', s1', ltr1', s2, ltr2, statA, sv1', sv2', statOO)) \wedge$ 
   $lltrv2 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$ 
   $lappend (l\list-of trv2) (lltrv2 (if trv1 = [] then L else R,$ 
   $w1', w2', s1', ltr1', s2, ltr2, statA, sv1', sv2', statOO))$ 

```

```

proof–
show ?thesis
using h-t-lnever-L[OF assms] apply(elim exE)
subgoal for w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1']) apply(rule exI[of - ltr1'])
apply(rule exI[of - trv1]) apply(rule exI[of - sv1'])
apply(rule exI[of - trv2]) apply(rule exI[of - sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal by simp
  subgoal by simp
  subgoal by simp
  subgoal by simp

```

subgoal by simp
subgoal by simp
subgoal unfolding ltrv1-def apply(subst f-length-ge1[OF len])
unfolding h-def t-def by auto
subgoal unfolding ltrv2-def apply(subst f-length-ge1[OF len])
unfolding h-def t-def by auto . .
qed

lemma h-t-not-lnever-L:
assumes unvw: unwindCond Δ
and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$
and r : reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and l': lnever isIntO ltr1 \neg isIntO s2
and len: llength ltr1 > Suc 0 llength ltr2 > Suc 0
and l: trn = L \neg lnever isSecO ltr1
shows $\exists w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO.$
 $\chi\chi s1 ltr1 tr1 s1' s1'' ltr1' \wedge$
 $\chi\chi^3 \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$
 \wedge
 $h-t (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $((trv1, sv1'', trv2, sv2'', statA, statOO),$
 $(R, w1', w2', s1'', s1'' \$ ltr1', s2, ltr2, statA, sv1'', sv2'', statOO))$
proof –
have s1: \neg isIntO s1 using l' ltr1
by (metis Opt.lvalidFromS-def l(2) lhd-LCons llist.exhaust llist.pred-inject(1)
llist.pred-inject(2))

obtain tr1 s1' s1'' ltr1'
where $\chi\chi$: $\chi\chi s1 ltr1 tr1 s1' s1'' ltr1'$
using isSecO- $\chi\chi$ [OF ltr1 l'(1) l(2)] by auto

define tp where
 $tp = (SOME k. case k of (tr1, s1', s1'', ltr1') \Rightarrow$
 $\chi\chi s1 ltr1 tr1 s1' s1'' ltr1')$

have 0: case tp of (tr1, s1', s1'', ltr1') \Rightarrow
 $\chi\chi s1 ltr1 tr1 s1' s1'' ltr1'$
using $\chi\chi$ unfolding tp-def apply– apply(rule someI-ex)
apply(rule exI[of - (tr1, s1', s1'', ltr1')]) by auto

obtain tr1 s1' s1'' ltr1' where
 $tp: tp = (tr1, s1', s1'', ltr1') \text{ by(cases tp, auto)}$

have $\chi\chi$: $\chi\chi s1 ltr1 tr1 s1' s1'' ltr1'$
using 0 unfolding tp by auto

obtain $w1' w2' trv1 sv1'' trv2 sv2'' statOO$
where $\chi^3: \chi^3' \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$
using $unwindCond-ex-\chi^3'[OF unw \Delta r, of tr1 s1' s1'']$
using $\chi\chi l' s1$ **unfolding** $\chi\chi-def$ **by** $auto$

define tp' **where**
 $tp' = (SOME k'. case k' of (w1',w2',trv1,sv1'',trv2,sv2'',statOO) \Rightarrow$
 $\chi^3' \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2 sv2''$
 $statOO)$

have $1: case tp' of (w1',w2',trv1,sv1'',trv2,sv2'',statOO) \Rightarrow$
 $\chi^3' \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2 sv2''$
 $statOO$

using χ^3' **unfolding** $tp'-def$ **apply**– **apply**($rule someI-ex$)
apply($rule exI[of - (w1',w2',trv1,sv1'',trv2,sv2'',statOO)]$) **by** $auto$

obtain $w1' w2' trv1 sv1'' trv2 sv2'' statOO$ **where**
 $tp': tp' = (w1',w2',trv1,sv1'',trv2,sv2'',statOO)$ **by**($cases tp', auto$)

have $\chi^3: \chi^3' \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$
using 1 **unfolding** tp' **by** $auto$

show $?thesis$
apply($rule exI[of - w1']$) **apply**($rule exI[of - w2']$)
apply($rule exI[of - tr1]$) **apply**($rule exI[of - s1']$) **apply**($rule exI[of - s1'']$)
apply($rule exI[of - ltr1]$)
apply($rule exI[of - trv1]$) **apply**($rule exI[of - sv1'']$) **apply**($rule exI[of - trv2]$)
apply($rule exI[of - sv2'']$)
apply($rule exI[of - statOO]$)
apply($intro conjI$)
subgoal **using** $\chi\chi$.
subgoal **using** χ^3' .
subgoal **using** l **unfolding** $h-t.simps$
unfolding $tp-def[symmetric]$ tp **apply** $simp$
unfolding $tp'-def[symmetric]$ tp' **by** $simp$.
qed

lemma $lltrv1-lltrv2-not-lnever-L:$
assumes $unw: unwindCond \Delta$
and $\Delta: \Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$
and $r: reachO s1 reachO s2 reachV sv1 reachV sv2$
and $ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1$
and $l': lnever isIntO ltr1 \neg isIntO s2$
and $len: llength ltr1 > Suc 0 llength ltr2 > Suc 0$
and $l: trn = L \neg lnever isSecO ltr1$

shows $\exists w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO.$
 $\chi\chi s1 ltr1 tr1 s1' s1'' ltr1' \wedge$
 $\chi3' \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$
 \wedge
 $lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $lappend (l\text{list-of } trv1) (lltrv1 (R, w1', w2', s1'', s1'' \$ ltr1', s2, ltr2, statA, sv1'', sv2'', statOO))$
 \wedge
 $lltrv2 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $lappend (l\text{list-of } trv2) (lltrv2 (R, w1', w2', s1'', s1'' \$ ltr1', s2, ltr2, statA, sv1'', sv2'', statOO))$

proof –

show *?thesis*
using *h-t-not-lnever-L[OF assms]* **apply** (*elim exE*)
subgoal for $w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO$
apply (*rule exI[of - w1']*) **apply** (*rule exI[of - w2']*)
apply (*rule exI[of - tr1]*) **apply** (*rule exI[of - s1']*) **apply** (*rule exI[of - s1'']*)
apply (*rule exI[of - ltr1']*)
apply (*rule exI[of - trv1]*) **apply** (*rule exI[of - sv1'']*) **apply** (*rule exI[of - trv2]*)
apply (*rule exI[of - sv2'']*)
apply (*rule exI[of - statOO]*)
apply (*intro conjI*)
subgoal by *simp*
subgoal by *simp*
subgoal unfolding *lltrv1-def* **apply** (*subst f-length-ge1[OF len]*)
unfolding *h-def t-def* **by** *simp*
subgoal unfolding *lltrv2-def* **apply** (*subst f-length-ge1[OF len]*)
unfolding *h-def t-def* **by** *simp . .*

qed

lemma *h-t-lnever-R*:

assumes *unw: unwindCond* Δ
and $\Delta: \Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$
and *r: reachO s1 reachO s2 reachV sv1 reachV sv2*
and *ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2*
and *l': $\neg isIntO s1 lnever isIntO ltr2$*
and *len: llength ltr1 > Suc 0 llength ltr2 > Suc 0*
and *l: trn = R lnever isSecO ltr2*
shows $\exists w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO.$
 $ltr2 = s2 \$ ltr2' \wedge validTransO (s2, s2') \wedge$
 $Opt.lvalidFromS s2' ltr2' \wedge Opt.lcompletedFrom s2' ltr2' \wedge lnever isIntO ltr2' \wedge$

$\omega4 \Delta w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2' statOO \wedge$
 $h-t (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $((trv1, sv1', trv2, sv2', statA, statOO),$
 $(if trv2 = [] then R else L,$
 $w1', w2', s1, ltr1, s2', ltr2', statA, sv1', sv2', statOO))$

proof –

have $s2: \neg \text{isIntO } s2$ **using** $l' \text{ ltr2}$
by (*metis* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-def* *lfinite-LNil* *lhd-LCons* *llist.exhaust* *llist.pred-inject*(2))

obtain $\text{ltr2}'$ **where** $\text{ltr2}_4: \text{ltr2} = s2 \ \$ \ \text{ltr2}'$
by (*metis* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-def* *lfinite-LNil* *llist.exhaust-sel* *ltr2*(1) *ltr2*(2))
hence $\text{ltr2}': \text{ltr2}' = \text{ltl } \text{ltr2}$ **by** *auto*
have $\text{ltr2}'\text{ne}: \text{ltr2}' \neq []$ **using** $\text{len}(2)$ **unfolding** ltr2_4
by (*metis* *One-nat-def* *llength-LCons* *llength-LNil* *one-eSuc* *one-enat-def* *order-less-irrefl*)
define $s2'$ **where** $s2': s2' = \text{lhd } (\text{ltl } \text{ltr2})$
have $v4: \text{validTransO } (s2, s2')$ **and** $vv4: \text{Opt.lvalidFromS } s2' \ \text{ltr2}' \ \text{Opt.lcompletedFrom } s2' \ \text{ltr2}'$
using $\text{ltr2 } \text{ltr2}'\text{ne}$ **unfolding** $\text{ltr2}_4 \ s2'$
by (*metis* *Opt.lcompletedFrom-LCons* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-Cons-iff* $\text{ltr2}' \ \text{ltr2}_4$)+

have $\text{is2}': \neg \text{isIntO } s2'$ **and** $\text{lnever } \text{isIntO } \text{ltr2}'$
using $l'(2)$ **unfolding** ltr2_4
by (*metis* *llist.exhaust-sel* *llist.pred-inject*(2) $\text{ltr2}' \ \text{ltr2}'\text{ne } s2'$)+

have $\text{iss2}: \neg \text{isSecO } s2$
using $l(2) \ \text{ltr2}_4$ **by** *auto*

obtain $w1' \ w2' \ \text{trv1} \ \text{sv1}' \ \text{trv2} \ \text{sv2}' \ \text{statOO}$
where $\omega_4: \omega_4 \ \Delta \ w1 \ w2 \ w1' \ w2' \ s1 \ s2 \ (\text{lhd } (\text{ltl } \text{ltr2})) \ \text{statA} \ \text{sv1} \ \text{trv1} \ \text{sv1}' \ \text{sv2} \ \text{trv2} \ \text{sv2}' \ \text{statOO}$
using *unwindCond-ex* ω_4 [*OF* *unw* $\Delta \ r \ l'(1) \ v4 \ s2 \ \text{is2}' \ \text{iss2}$] $s2'$ **by** *auto*

define tp' **where**
 $tp' = (\text{SOME } k'. \text{ case } k' \text{ of } (w1', w2', \text{trv1}, \text{sv1}', \text{trv2}, \text{sv2}', \text{statOO}) \Rightarrow$
 $\omega_4 \ \Delta \ w1 \ w2 \ w1' \ w2' \ s1 \ s2 \ (\text{lhd } (\text{ltl } \text{ltr2})) \ \text{statA} \ \text{sv1} \ \text{trv1} \ \text{sv1}' \ \text{sv2} \ \text{trv2} \ \text{sv2}' \ \text{statOO})$

have $1: \text{ case } tp' \text{ of } (w1', w2', \text{trv1}, \text{sv1}', \text{trv2}, \text{sv2}', \text{statOO}) \Rightarrow$
 $\omega_4 \ \Delta \ w1 \ w2 \ w1' \ w2' \ s1 \ s2 \ s2' \ \text{statA} \ \text{sv1} \ \text{trv1} \ \text{sv1}' \ \text{sv2} \ \text{trv2} \ \text{sv2}' \ \text{statOO}$
using ω_4 **unfolding** $tp'\text{-def}$ $s2'$ **apply**– **apply**(*rule* *someI-ex*)
apply(*rule* *exI*[*of* - $(w1', w2', \text{trv1}, \text{sv1}', \text{trv2}, \text{sv2}', \text{statOO})$]) **by** *auto*

obtain $w1' \ w2' \ \text{trv1} \ \text{sv1}' \ \text{trv2} \ \text{sv2}' \ \text{statOO}$ **where**
 $tp': tp' = (w1', w2', \text{trv1}, \text{sv1}', \text{trv2}, \text{sv2}', \text{statOO})$ **by**(*cases* tp' , *auto*)

have $\omega_4: \omega_4 \ \Delta \ w1 \ w2 \ w1' \ w2' \ s1 \ s2 \ s2' \ \text{statA} \ \text{sv1} \ \text{trv1} \ \text{sv1}' \ \text{sv2} \ \text{trv2} \ \text{sv2}' \ \text{statOO}$
using 1 **unfolding** tp' **by** *auto*

```

show ?thesis
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])
apply(rule exI[of - trv1]) apply(rule exI[of - sv1'])
apply(rule exI[of - trv2]) apply(rule exI[of - sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal using len l unfolding h-t.simps apply simp
  unfolding tp'-def[symmetric] tp' s2' ltr2' by simp .
qed

```

```

lemma lltrv1-lltrv2-lnever-R:
assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and l':  $\neg$  isIntO s1 lnever isIntO ltr2
and len: llength ltr1 > Suc 0 llength ltr2 > Suc 0
and l: trn = R lnever isSecO ltr2
shows  $\exists w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO.$ 
  ltr2 = s2 $ ltr2'  $\wedge$  validTransO (s2,s2')  $\wedge$ 
  Opt.lvalidFromS s2' ltr2'  $\wedge$  Opt.lcompletedFrom s2' ltr2'  $\wedge$  lnever isIntO ltr2'  $\wedge$ 

 $\omega_4 \Delta w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2' statOO \wedge$ 
  lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) =
    lappend (llist-of trv1) (lltrv1 (if trv2 = [] then R else L,
      w1',w2',s1,ltr1,s2',ltr2',statA,sv1',sv2',statOO))  $\wedge$ 
  lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) =
    lappend (llist-of trv2) (lltrv2 (if trv2 = [] then R else L,
      w1',w2',s1,ltr1,s2',ltr2',statA,sv1',sv2',statOO))

```

```

proof –
show ?thesis
using h-t-lnever-R[OF assms] apply(elim exE)
subgoal for w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])
apply(rule exI[of - trv1]) apply(rule exI[of - sv1'])
apply(rule exI[of - trv2]) apply(rule exI[of - sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal by simp
  subgoal by simp
  subgoal by simp

```

subgoal by simp
subgoal by simp
subgoal by simp
subgoal unfolding ltrv1-def apply(subst f-length-ge1[OF len])
unfolding h-def t-def by auto
subgoal unfolding ltrv2-def apply(subst f-length-ge1[OF len])
unfolding h-def t-def by auto . .
qed

lemma h-t-not-lnever-R:
assumes unvw: unwindCond Δ
and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and l': $\neg isIntO s1 lnever isIntO ltr2$
and len: llength ltr1 > Suc 0 llength ltr2 > Suc 0
and l: trn = R $\neg lnever isSecO ltr2$
shows $\exists w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO.$
 $\chi\chi s2 ltr2 tr2 s2' s2'' ltr2' \wedge$
 $\chi\chi' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$
 \wedge
 $h-t (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $((trv1, sv1'', trv2, sv2'', statA, statOO),$
 $(L, w1', w2', s1, ltr1, s2'', s2'' \$ ltr2', statA, sv1'', sv2'', statOO))$
proof –
have s2: $\neg isIntO s2$ using l' ltr2
by (metis Simple-Transition-System.lvalidFromS-def l(2) lhd-LCons llist.pred-inject(1)
 $l\text{list.pred-inject}(2) \text{neq-LNil-conv})$
obtain tr2 s2' s2'' ltr2'
where $\chi\chi$: $\chi\chi s2 ltr2 tr2 s2' s2'' ltr2'$
using isSecO- $\chi\chi$ [OF ltr2 l'(2) l(2)] by auto
define tp where
 $tp = (SOME k. \text{case } k \text{ of } (tr2, s2', s2'', ltr2') \Rightarrow$
 $\chi\chi s2 ltr2 tr2 s2' s2'' ltr2')$
have 0: case tp of (tr2, s2', s2'', ltr2') \Rightarrow
 $\chi\chi s2 ltr2 tr2 s2' s2'' ltr2'$
using $\chi\chi$ unfolding tp-def apply- apply(rule someI-ex)
apply(rule exI[of - (tr2, s2', s2'', ltr2')]) by auto
obtain tr2 s2' s2'' ltr2' where
 $tp: tp = (tr2, s2', s2'', ltr2')$ **by(cases tp, auto)**
have $\chi\chi$: $\chi\chi s2 ltr2 tr2 s2' s2'' ltr2'$


```

using 0 unfolding tp by auto

obtain w1' w2' trv1 sv1'' trv2 sv2'' statOO
where  $\chi_4'$ :  $\chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$ 
using unwindCond-ex- $\chi_4'$ [OF unW  $\Delta$  r, of tr2 s2' s2'']
using  $\chi\chi$  l' s2 unfolding  $\chi\chi$ -def by auto

define tp' where
tp' = (SOME k'. case k' of (w1',w2',trv1,sv1'',trv2,sv2'',statOO)  $\Rightarrow$ 
 $\chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$ )

have 1: case tp' of (w1',w2',trv1,sv1'',trv2,sv2'',statOO)  $\Rightarrow$ 
 $\chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$ 
using  $\chi_4'$  unfolding tp'-def apply- apply(rule someI-ex)
apply(rule exI[of - (w1',w2',trv1,sv1'',trv2,sv2'',statOO)]) by auto

obtain w1' w2' trv1 sv1'' trv2 sv2'' statOO where
tp': tp' = (w1',w2',trv1,sv1'',trv2,sv2'',statOO) by(cases tp', auto)

have  $\chi_4'$ :  $\chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$ 
using 1 unfolding tp' by auto

show ?thesis
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - tr2]) apply(rule exI[of - s2']) apply(rule exI[of - s2''])
apply(rule exI[of - ltr2'])
apply(rule exI[of - trv1]) apply(rule exI[of - sv1'']) apply(rule exI[of - trv2])
apply(rule exI[of - sv2''])
apply(rule exI[of - statOO])
apply(intro conjI)
subgoal using  $\chi\chi$  .
subgoal using  $\chi_4'$  .
subgoal using l unfolding h-t.simps
unfolding tp-def[symmetric] tp apply simp
unfolding tp'-def[symmetric] tp' by auto .
qed

lemma lltrv1-lltrv2-not-lnever-R:
assumes unW: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and l':  $\neg isIntO s1 lnever isIntO ltr2$ 

```

and $len: llength\ ltr1 > Suc\ 0\ llength\ ltr2 > Suc\ 0$
and $l: trn = R \ \neg\ lnever\ isSecO\ ltr2$
shows $\exists\ w1'\ w2'\ tr2\ s2'\ s2''\ ltr2'\ trv1\ sv1''\ trv2\ sv2''\ statOO.$
 $\chi\chi\ s2\ ltr2\ tr2\ s2'\ s2''\ ltr2' \wedge$
 $\chi4'\ \Delta\ w1\ w2\ w1'\ w2'\ s1\ s2\ tr2\ s2'\ s2''\ statA\ sv1\ trv1\ sv1''\ sv2\ trv2\ sv2''\ statOO$
 \wedge
 $lltrv1\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $lappend\ (l\text{list-of}\ trv1)\ (lltrv1\ (L, w1', w2', s1, ltr1, s2'', s2''\ \$\ ltr2', statA, sv1'', sv2'', statOO))$
 \wedge
 $lltrv2\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $lappend\ (l\text{list-of}\ trv2)\ (lltrv2\ (L, w1', w2', s1, ltr1, s2'', s2''\ \$\ ltr2', statA, sv1'', sv2'', statOO))$
proof –
show *?thesis*
using $h\text{-}t\text{-not-lnever-R}[OF\ assms]$ **apply** $(elim\ exE)$
subgoal for $w1'\ w2'\ tr2\ s2'\ s2''\ ltr2'\ trv1\ sv1''\ trv2\ sv2''\ statOO$
apply $(rule\ exI[of\ -\ w1'])$ **apply** $(rule\ exI[of\ -\ w2'])$
apply $(rule\ exI[of\ -\ tr2])$ **apply** $(rule\ exI[of\ -\ s2'])$ **apply** $(rule\ exI[of\ -\ s2''])$
apply $(rule\ exI[of\ -\ ltr2'])$
apply $(rule\ exI[of\ -\ trv1])$ **apply** $(rule\ exI[of\ -\ sv1''])$ **apply** $(rule\ exI[of\ -\ trv2])$
apply $(rule\ exI[of\ -\ sv2''])$
apply $(rule\ exI[of\ -\ statOO])$
apply $(intro\ conjI)$
subgoal by *simp*
subgoal by *simp*
subgoal unfolding $lltrv1\text{-}def$ **apply** $(subst\ f\text{-}length\text{-}ge1[OF\ len])$
unfolding $h\text{-}def\ t\text{-}def$ **by** *simp*
subgoal unfolding $lltrv2\text{-}def$ **apply** $(subst\ f\text{-}length\text{-}ge1[OF\ len])$
unfolding $h\text{-}def\ t\text{-}def$ **by** *simp* . .
qed

lemma $f\text{-}not\text{-}LNil: ltr1 \neq [] \vee ltr2 \neq [] \implies$
 $f\ (w1, w2, trn, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \neq []$
apply $(cases\ llength\ ltr1 \leq Suc\ 0 \vee llength\ ltr2 \leq Suc\ 0)$
subgoal apply $(subst\ f\text{-}length\text{-}1)$ **by** *auto*
subgoal apply $(subst\ f\text{-}length\text{-}ge1)$ **by** *auto* .

lemma $lvalidFromS\ lltrv1:$
assumes $unw: unwindCond\ \Delta$
and $\Delta: \Delta\ \infty\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
and $r: reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$
and $ltr1: Opt.lvalidFromS\ s1\ ltr1\ Opt.lcompletedFrom\ s1\ ltr1\ lnever\ isIntO\ ltr1$
and $ltr2: Opt.lvalidFromS\ s2\ ltr2\ Opt.lcompletedFrom\ s2\ ltr2\ lnever\ isIntO\ ltr2$
shows $Van.lvalidFromS\ sv1\ (lltrv1\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO))$
proof –
{fix $n\ sv1\ lltrv1$
assume $\exists\ trn\ w1\ w2\ s1\ ltr1\ s2\ ltr2\ statA\ sv2\ statO.$

```

n = w1 ∧
ltrv1 = lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ∧
Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧ lnever isIntO ltr1 ∧

  Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧ lnever isIntO ltr2
hence Van.llvalidFromS n sv1 ltrv1
proof(coinduct rule: Van.llvalidFromS.coinduct[of λn sv1 ltrv1.
  ∃ trn w1 w2 s1 ltr1 s2 ltr2 statA sv2 statO.
    n = w1 ∧
    ltrv1 = lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ∧
    Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
    reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
    Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧ lnever isIntO ltr1 ∧

      Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧ lnever isIntO ltr2])
case (llvalidFromS n sv1 ltrv1)
then obtain trn w1 w2 s1 ltr1 s2 ltr2 statA sv2 statO
where n: n = w1 and
ltrv1: ltrv1 = lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
by auto
have isi3: ¬ isIntO s1 using ltr1
by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
llist.pred-inject(2))
have isi4: ¬ isIntO s2 using ltr2
by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
llist.pred-inject(2))

show ?case proof(cases ltr1 = [[]] ∧ ltr2 = [[]])
  case True note ltr14 = True
  hence ltrv1: ltrv1 = [[]] unfolding ltrv1 by simp
  show ?thesis unfolding ltrv1 apply(rule Van.llvalidFromS-selectLNil) by
auto
next
case False hence ltr14: ltr1 ≠ [[]] ∨ ltr2 ≠ [[]] by auto
show ?thesis proof(cases llength ltr1 ≤ Suc 0 ∨ llength ltr2 ≤ Suc 0)
  case True note ltr14 = ltr14 True
  hence ltrv1: ltrv1 = [[sv1]] unfolding ltrv1 by simp
  show ?thesis unfolding ltrv1 apply(rule Van.llvalidFromS-selectSingl) by
auto
next
case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0
by auto
show ?thesis proof(cases trn)

```

```

case L note trn = L note current = current L
show ?thesis
proof(cases lnever isSecO ltr1)
  case True note current = current True
  obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
    ωω: ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
    lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
    ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
  and trn': trn' = (if trv1 = [] then L else R)
  and ltrv1: ltrv1 =
    lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-L[OF unω Δ r ltr1 isi4 current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1' ≡ lltrv1 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
  have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..

show ?thesis
proof(cases trv1 = [])
  case True note trv1 = True
  have sv1': sv1' = sv1
  using ω3 unfolding ω3-def by (simp add: trv1)
  show ?thesis
  apply(rule Van.llvalidFromS-selectDelay)
  apply(rule exI[of - w1']) apply(rule exI[of - n])
  apply(rule exI[of - sv1]) apply(rule exI[of - ltrv1'])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1 trv1 by simp
  subgoal using ω3 unfolding ω3-def trv1 n by simp
  subgoal apply(rule disjI1)
    apply(rule exI[of - trn'])
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1']) apply(rule exI[of - ltr1'])
    apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
    apply(rule exI[of - statA]) apply(rule exI[of - sv2']) apply(rule
exI[of - statOO])
  apply(intro conjI)
  subgoal ..
  subgoal unfolding ltrv1' trn' trv1 sv1' using trn by simp
  subgoal using ω3 unfolding ω3-def sv1' by simp
  subgoal using ω3 unfolding ω3-def
  by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
  subgoal by fact subgoal by fact
  subgoal using ω3 unfolding ω3-def
  by (metis Nil-is-append-conv Van.reach-validFromS-reach last-snoc

```

```

not-Cons-self2 r(4))
  subgoal using  $\omega\omega$  by simp subgoal using  $\omega\omega$  by simp subgoal
using  $\omega\omega$  by simp
  subgoal by fact subgoal by fact subgoal by fact . .
next
case False note trv1 = False
show ?thesis
apply(rule Van.llvalidFromS-selectlappend)
apply(rule exI[of - sv1]) apply(rule exI[of - trv1])
apply(rule exI[of - sv1']) apply(rule exI[of - w1'])
apply(rule exI[of - ltrv1']) apply(rule exI[of - n])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv1 .
  subgoal using  $\omega\mathfrak{3}$  unfolding  $\omega\mathfrak{3}$ -def
by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
  subgoal by fact
  subgoal using  $\omega\mathfrak{3}$  unfolding  $\omega\mathfrak{3}$ -def
  by (metis Van.validFromS-def Van.validS-validTrans list.sel(1)
not-Cons-self2 snoc-eq-iff-butlast trv1)
  subgoal apply(rule disjI1)
  apply(rule exI[of - trn'])
  apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule exI[of
- s1']) apply(rule exI[of - ltr1'])
  apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
  apply(rule exI[of - statA]) apply(rule exI[of - sv2']) apply(rule
exI[of - statOO])
  apply(intro conjI)
  subgoal ..
  subgoal using trv1 unfolding ltrv1' trn' by auto
  subgoal using  $\omega\mathfrak{3}$  unfolding  $\omega\mathfrak{3}$ -def by simp
  subgoal using  $\omega\mathfrak{3}$  unfolding  $\omega\mathfrak{3}$ -def
  by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(1) snd-conv)
  subgoal by fact
  subgoal using  $\omega\mathfrak{3}$  unfolding  $\omega\mathfrak{3}$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\omega\mathfrak{3}$  unfolding  $\omega\mathfrak{3}$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$ 
  using llist-all-lappend-llist-of ltr1(3) by blast
  subgoal using  $\omega\omega$  using ltr2(1) by fastforce
  subgoal by fact
  subgoal by fact . .
qed

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next
  case False note current = current False
  obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
   $\chi\chi$ :  $\chi\chi$  s1 ltr1 tr1 s1' s1'' ltr1' and
   $\chi\mathcal{I}$ ':  $\chi\mathcal{I}$ '  $\Delta$  w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
  trv2 sv2'' statOO
  and ltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

  using lltrv1-lltrv2-not-lnever-L[OF unW  $\Delta$  r ltr1 isi4 current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (R,w1',w2',s1'',s1'' $
  ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
  have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..

show ?thesis apply(rule Van.llvalidFromS-selectlappend)
apply(rule exI[of - sv1]) apply(rule exI[of - trv1])
apply(rule exI[of - sv1'']) apply(rule exI[of - w1'])
apply(rule exI[of - ltrv1']) apply(rule exI[of - w1])
apply(intro conjI)
  subgoal unfolding n .. subgoal ..
  subgoal using ltrv1 .
  subgoal using  $\chi\mathcal{I}$ ' unfolding  $\chi\mathcal{I}$ '-def
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
  hd-append2)
  subgoal using  $\chi\mathcal{I}$ ' unfolding  $\chi\mathcal{I}$ '-def by simp
  subgoal using  $\chi\mathcal{I}$ ' unfolding  $\chi\mathcal{I}$ '-def
  by (metis Van.validFromS Van.validS-validTrans Simple-Transition-System.validFromS-def

  append-is-Nil-conv not-Cons-self2)
  subgoal apply(rule disjI1)
  apply(rule exI[of - R])
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1'']) apply(rule exI[of - s1'' $ ltr1'])
  apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
  apply(rule exI[of - statA]) apply(rule exI[of - sv2'']) apply(rule exI[of
  - statOO])
  apply(intro conjI)
  subgoal ..
  subgoal unfolding ltrv1' ..
  subgoal using  $\chi\mathcal{I}$ ' unfolding  $\chi\mathcal{I}$ '-def by simp
  subgoal using  $\chi\mathcal{I}$ ' unfolding  $\chi\mathcal{I}$ '-def
  by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$   $\chi\chi$ -def

  append-is-Nil-conv last-snoc not-Cons-self2 r(1))
  subgoal by fact
  subgoal using  $\chi\mathcal{I}$ ' unfolding  $\chi\mathcal{I}$ '-def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3))

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snoc-eq-iff-butlast)
  subgoal using  $\chi\beta'$  unfolding  $\chi\beta'$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach r(4))
snoc-eq-iff-butlast)
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
  using llist-all-lappend-llist-of ltr1(3) by blast
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def using ltr2(1) by fastforce
  subgoal by fact
  subgoal by fact ..
qed
next
case R note trn = R note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
  case True note current = current True
  obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
   $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
  lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
   $\omega\omega$ :  $\omega\omega \Delta w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'$ 
statOO
  and trn': trn' = (if trv2 = [] then R else L)
  and ltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-R[OF un $\omega \Delta r$  ltr2(1,2) isi3 ltr2(3) current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
  have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..

show ?thesis
proof(cases trv1 = [])
  case True note trv1 = True
  have sv1': sv1' = sv1
  using  $\omega\omega$  unfolding  $\omega\omega$ -def by (simp add: trv1)
  show ?thesis
  apply(rule Van.llvalidFromS-selectDelay)
  apply(rule exI[of - w1']) apply(rule exI[of - n])
  apply(rule exI[of - sv1]) apply(rule exI[of - ltrv1'])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1 trv1 by simp
  subgoal using  $\omega\omega$  unfolding  $\omega\omega$ -def trv1 n by simp
  subgoal apply(rule disjI1)
    apply(rule exI[of - trn'])
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])

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    apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
    apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
    apply(rule exI[of - statA]) apply(rule exI[of - sv2]) apply(rule
exI[of - statOO])
    apply(intro conjI)
    subgoal ..
    subgoal unfolding ltrv1' trn' trv1 sv1' using trn by simp
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def sv1' by simp
    subgoal by fact
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(2) snd-conv)
    subgoal by fact
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Nil-is-append-conv Van.reach-validFromS-reach last-snoc
not-Cons-self2 r(4))
    subgoal by fact subgoal by fact subgoal by fact
    subgoal using  $\omega\omega$  by simp subgoal using  $\omega\omega$  by simp subgoal
using  $\omega\omega$  by simp . .
next
case False note trv1 = False
show ?thesis
apply(rule Van.llvalidFromS-selectlappend)
apply(rule exI[of - sv1]) apply(rule exI[of - trv1])
apply(rule exI[of - sv1]) apply(rule exI[of - w1])
apply(rule exI[of - ltrv1]) apply(rule exI[of - n])
apply(intro conjI)
    subgoal .. subgoal ..
    subgoal using ltrv1 .
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
    subgoal by fact
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Van.validFromS-def Van.validS-validTrans list.sel(1)
not-Cons-self2 snoc-eq-iff-butlast trv1)
    subgoal apply(rule disjI1)
    apply(rule exI[of - trn])
    apply(rule exI[of - w1]) apply(rule exI[of - w2]) apply(rule exI[of
- s1]) apply(rule exI[of - ltr1])
    apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
    apply(rule exI[of - statA]) apply(rule exI[of - sv2]) apply(rule
exI[of - statOO])
    apply(intro conjI)
    subgoal ..
    subgoal using trv1 unfolding ltrv1' trn' by auto
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
    subgoal by fact
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(2) snd-conv)

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      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
      subgoal by fact subgoal by fact subgoal by fact
      subgoal using  $\omega\omega$  by auto
      subgoal using  $\omega\omega$  by auto
      subgoal using  $\omega\omega$ 
      using llist-all-lappend-llist-of ltr1(3) by blast . .
qed
next
case False note current = current False
obtain  $w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO$  where
 $\chi\chi: \chi\chi s2 ltr2 tr2 s2' s2'' ltr2'$  and
 $\chi_4': \chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2$ 
 $trv2 sv2'' statOO$ 
and  $ltrv1: ltrv1 =$ 
  lappend (llist-of trv1) (lltrv1 (L, w1', w2', s1, ltr1, s2'', s2'' $ ltr2',
statA, sv1'', sv2'', statOO))
  using lltrv1-lltrv2-not-lnever-R[OF un $\omega$   $\Delta$  r ltr2(1,2) isi3 ltr2(3)
current]
  unfolding ltrv1 by blast
  define  $ltrv1'$  where  $ltrv1': ltrv1' \equiv lltrv1 (L, w1', w2', s1, ltr1, s2'',$ 
 $s2'' $ ltr2', statA, sv1'', sv2'', statOO)$ 
  have  $ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'$ 
  unfolding ltrv1 ltrv1' ..

show ?thesis
proof(cases trv1 = [])
case True note trv1 = True
hence  $sv1'': sv1'' = sv1$ 
by (metis  $\chi_4'$ -def Simple-Transition-System.validFromS-Cons-iff  $\chi_4'$ 
append.simps(1))
have  $w1' < w1$  using trv1  $\chi_4'$  unfolding  $\chi_4'$ -def by auto
show ?thesis
apply(rule Van.llvalidFromS-selectDelay)
apply(rule exI[of - w1']) apply(rule exI[of - n])
apply(rule exI[of - sv1]) apply(rule exI[of - ltrv1])
apply(intro conjI)
subgoal ..
subgoal .. subgoal .. subgoal unfolding n by fact
subgoal apply(rule disjI1)
apply(rule exI[of - L])
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
apply(rule exI[of - statA]) apply(rule exI[of - sv2'']) apply(rule

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exI[of - statOO])
  apply(intro conjI)
  subgoal ..
  subgoal unfolding ltrv1 ltrv1' trv1 sv1'' by simp
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def sv1'' by simp
  subgoal by fact
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
  by (metis Opt.reach-validFromS-reach Nil-is-append-conv last-snoc
not-Cons-self2 r(2))
  subgoal by fact
  subgoal using  $\chi_4'$  r(4) unfolding  $\chi_4'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal by fact subgoal by fact subgoal by fact
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
    using llist-all-lappend-llist-of ltr2(3) by blast . .
next
case False note trv1 = False
show ?thesis
apply(rule Van.llvalidFromS-selectlappend)
apply(rule exI[of - sv1]) apply(rule exI[of - trv1])
apply(rule exI[of - sv1''])
apply(rule exI[of - w1'])
apply(rule exI[of - ltrv1'])
apply(rule exI[of - n])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv1 .
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
  subgoal using trv1 .
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
  by (metis Simple-Transition-System.validFromS-def Van.validS-validTrans
list.sel(1)
  not-Cons-self2 snoc-eq-iff-butlast trv1)
  subgoal apply(rule disjI1)
  apply(rule exI[of - L])
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
  apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
  apply(rule exI[of - statA]) apply(rule exI[of - sv2'']) apply(rule
exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal unfolding ltrv1' ..
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
  subgoal by fact
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def

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    by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$ 
 $\chi\chi$ -def
    append-is-Nil-conv last-snoc not-Cons-self2 r(2))
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
    using llist-all-lappend-llist-of ltr2(3) by blast . .
    qed
  qed
  qed
  qed
  qed
  qed
}
thus ?thesis apply—apply(rule Van.llvalidFromS-imp-lvalidFromS)
using assms by blast
qed

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lemma lvalidFromS-lltrv2:
  assumes unw: unwindCond  $\Delta$ 
  and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
  shows Van.lvalidFromS sv2 (lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
proof—
  {fix n sv2 ltrv2
  assume  $\exists trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 statO.$ 
     $n = w2 \wedge$ 
     $ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) \wedge$ 
     $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$ 
     $reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$ 
     $Opt.lvalidFromS s1 ltr1 \wedge Opt.lcompletedFrom s1 ltr1 \wedge lnever isIntO ltr1 \wedge$ 
     $Opt.lvalidFromS s2 ltr2 \wedge Opt.lcompletedFrom s2 ltr2 \wedge lnever isIntO ltr2$ 
  hence Van.lvalidFromS n sv2 ltrv2
  proof(coinduct rule: Van.llvalidFromS.coinduct[of  $\lambda n sv2 ltrv2.$ 
 $\exists trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 statO.$ 
     $n = w2 \wedge$ 

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    ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ∧
    Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
    reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
    Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧ lnever isIntO ltr1 ∧

    Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧ lnever isIntO ltr2])
case (llvalidFromS n sv2 ltrv2)
then obtain trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 statO
where n: n = w2 and
ltrv2: ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
by auto
have isi3: ¬ isIntO s1 using ltr1
by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
llist.pred-inject(2))
have isi4: ¬ isIntO s2 using ltr2
by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
llist.pred-inject(2))

show ?case proof(cases ltr1 = [] ∧ ltr2 = [])
case True note ltr14 = True
hence ltrv2: ltrv2 = [] unfolding ltrv2 by simp
show ?thesis unfolding ltrv2 apply(rule Van.llvalidFromS-selectLNil) by
auto
next
case False hence ltr14: ltr1 ≠ [] ∨ ltr2 ≠ [] by auto
show ?thesis proof(cases llength ltr1 ≤ Suc 0 ∨ llength ltr2 ≤ Suc 0)
case True note ltr14 = ltr14 True
hence ltrv2: ltrv2 = [[sv2]] unfolding ltrv2 by simp
show ?thesis unfolding ltrv2 apply(rule Van.llvalidFromS-selectSingl) by
auto
next
case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0
by auto
show ?thesis proof(cases trn)
case L note trn = L note current = current L
show ?thesis
proof(cases lnever isSecO ltr1)
case True note current = current True
obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
ωω: ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
and trn': trn' = (if trv1 = [] then L else R)
and ltrv2: ltrv2 =

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      lappend (llist-of trv2) (lltrv2(trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
    using lltrv1-lltrv2-lnever-L[OF unv Δ r ltr1 isi4 current]
    unfolding ltrv2 by blast
    define ltrv2' where ltrv2': ltrv2' ≡ lltrv2 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
    have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
    unfolding ltrv2 ltrv2' ..

show ?thesis
proof(cases trv2 = [])
  case True note trv2 = True
  have sv2': sv2' = sv2
  using ω3 unfolding ω3-def by (simp add: trv2)
  show ?thesis
  apply(rule Van.llvalidFromS-selectDelay)
  apply(rule exI[of - w2']) apply(rule exI[of - n])
  apply(rule exI[of - sv2]) apply(rule exI[of - ltrv2'])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv2 trv2 by simp
  subgoal using ω3 unfolding ω3-def trv2 n by simp
  subgoal apply(rule disjI1)
    apply(rule exI[of - trn'])
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1']) apply(rule exI[of - ltr1'])
    apply(rule exI[of - s2]) apply(rule exI[of - ltr2'])
    apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule
exI[of - statOO])
    apply(intro conjI)
    subgoal ..
    subgoal unfolding ltrv2' trn' trv2 sv2' using trn by simp
    subgoal using ω3 unfolding ω3-def sv2' by simp
    subgoal using ω3 unfolding ω3-def
    by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using ω3 unfolding ω3-def
    by (metis Nil-is-append-conv Van.reach-validFromS-reach last-snoc
not-Cons-self2 r(3))
    subgoal by fact
    subgoal using ωω by simp subgoal using ωω by simp subgoal
using ωω by simp
    subgoal by fact subgoal by fact subgoal by fact . .
  next
  case False note trv2 = False
  show ?thesis
  apply(rule Van.llvalidFromS-selectlappend)
  apply(rule exI[of - sv2]) apply(rule exI[of - trv2])
  apply(rule exI[of - sv2']) apply(rule exI[of - w2'])

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apply(rule exI[of - ltrv2']) apply(rule exI[of - n])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv2 .
  subgoal using  $\omega^3$  unfolding  $\omega^3$ -def
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
    subgoal by fact
    subgoal using  $\omega^3$  unfolding  $\omega^3$ -def
    by (metis Van.validFromS-def Van.validS-validTrans append-is-Nil-conv
list.sel(1) not-Cons-self2 trv2)
      subgoal apply(rule disjI1)
      apply(rule exI[of - trn'])
      apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule exI[of
- s1']) apply(rule exI[of - ltr1'])
      apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])
      apply(rule exI[of - statA']) apply(rule exI[of - sv1']) apply(rule
exI[of - statOO'])
      apply(intro conjI)
      subgoal ..
      subgoal using trv2 unfolding ltrv2' trn' by auto
      subgoal using  $\omega^3$  unfolding  $\omega^3$ -def by simp
      subgoal using  $\omega^3$  unfolding  $\omega^3$ -def
      by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(1) snd-conv)
      subgoal by fact
      subgoal using  $\omega^3$  unfolding  $\omega^3$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
      subgoal using  $\omega^3$  unfolding  $\omega^3$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
      subgoal using  $\omega\omega$  by auto
      subgoal using  $\omega\omega$  by auto
      subgoal using  $\omega\omega$ 
      using l1ist-all-lappend-l1ist-of ltr1(3) by blast
      subgoal using  $\omega\omega$  using ltr2(1) by fastforce
      subgoal by fact
      subgoal by fact . .
    qed
  next
  case False note current = current False
  obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
   $\chi\chi$ :  $\chi\chi$  s1 ltr1 tr1 s1' s1'' ltr1' and
   $\chi^3$ :  $\chi^3$   $\Delta$  w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
  and ltrv2: ltrv2 =
lappend (l1ist-of trv2) (lltrv2 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

  using lltrv1-lltrv2-not-lnever-L[OF un $\omega$   $\Delta$  r ltr1 isi4 current]

```

```

unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2'  $\equiv$  ltrv2 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
  have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding ltrv2 ltrv2' ..

show ?thesis
proof(cases trv2 = [])
  case True note trv2 = True
  hence sv2'': sv2'' = sv2
  by (metis  $\chi^{3'}$ -def Simple-Transition-System.validFromS-Cons-iff  $\chi^{3'}$ 
append.simps(1))
  have w2' < w2 using trv2  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by auto
  show ?thesis
  apply(rule Van.llvalidFromS-selectDelay)
  apply(rule exI[of - w2  $\uparrow$ ]) apply(rule exI[of - n])
  apply(rule exI[of - sv2]) apply(rule exI[of - ltrv2])
  apply(intro conjI)
  subgoal ..
  subgoal .. subgoal .. subgoal unfolding n by fact
  subgoal apply(rule disjI1)
  apply(rule exI[of - R])
  apply(rule exI[of - w1  $\uparrow$ ]) apply(rule exI[of - w2  $\uparrow$ ])
  apply(rule exI[of - s1''  $\uparrow$ ]) apply(rule exI[of - s1'' $ ltr1  $\uparrow$ ])
  apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
  apply(rule exI[of - statA]) apply(rule exI[of - sv1'']) apply(rule
exI[of - statOO])
  apply(intro conjI)
  subgoal ..
  subgoal unfolding ltrv2 ltrv2' trv2 sv2'' by simp
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def sv2'' by simp
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
  by (metis Opt.reach-validFromS-reach Nil-is-append-conv last-snoc
not-Cons-self2 r(1))
  subgoal by fact
  subgoal using  $\chi^{3'}$  r(3) unfolding  $\chi^{3'}$ -def
  by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal by fact
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
  using llist-all-lappend-llist-of ltr1(3) by blast
  subgoal by fact subgoal by fact subgoal by fact . .
next
  case False note trv2 = False
  show ?thesis
  apply(rule Van.llvalidFromS-selectlappend)
  apply(rule exI[of - sv2]) apply(rule exI[of - trv2])
  apply(rule exI[of - sv2'']) apply(rule exI[of - w2  $\uparrow$ ])

```

```

apply(rule exI[of - ltrv2']) apply(rule exI[of - w2])
apply(intro conjI)
  subgoal unfolding n .. subgoal ..
  subgoal using ltrv2 .
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
  subgoal by fact
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
  by (metis Simple-Transition-System.validFromS-def Van.validS-validTrans
append-is-Nil-conv list.sel(1) not-Cons-self2 trv2)
  subgoal apply(rule disjI1)
  apply(rule exI[of - R])
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1']) apply(rule exI[of - s1'' $ ltr1'])
  apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])
  apply(rule exI[of - statA']) apply(rule exI[of - sv1']) apply(rule
exI[of - statOO'])
  apply(intro conjI)
  subgoal ..
  subgoal unfolding ltrv2' ..
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$  by simp
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
  by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$ 
XX-def
  append-is-Nil-conv last-snoc not-Cons-self2 r(1))
  subgoal by fact
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$ 
  using l1ist-all-lappend-l1ist-of ltr1(3) by blast
  subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  using ltr2(1) by fastforce
  subgoal by fact
  subgoal by fact . .
  qed
qed
next
case R note trn = R note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
  case True note current = current True
  obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
   $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'

```



```

lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
  ω4: ω4 Δ w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
and trn': trn' = (if trv2 = [] then R else L)
and ltrv2: ltrv2 =
  lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
using lltrv1-lltrv2-lnever-R[OF unω Δ r ltr2(1,2) isi3 ltr2(3) current]
  unfolding ltrv2 by blast
define ltrv2' where ltrv2': ltrv2' ≡ lltrv2 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding ltrv2 ltrv2' ..

show ?thesis
proof(cases trv2 = [])
  case True note trv2 = True
  have sv2': sv2' = sv2
  using ω4 unfolding ω4-def by (simp add: trv2)
  show ?thesis
  apply(rule Van.llvalidFromS-selectDelay)
  apply(rule exI[of - w2]) apply(rule exI[of - n])
  apply(rule exI[of - sv2]) apply(rule exI[of - ltrv2])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv2 trv2 by simp
  subgoal using ω4 unfolding ω4-def trv2 n by simp
  subgoal apply(rule disjI1)
    apply(rule exI[of - trn])
    apply(rule exI[of - w1]) apply(rule exI[of - w2])
    apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
    apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
    apply(rule exI[of - statA]) apply(rule exI[of - sv1]) apply(rule
exI[of - statOO])
  apply(intro conjI)
  subgoal ..
  subgoal unfolding ltrv2' trn' trv2 sv2' using trn by simp
  subgoal using ω4 unfolding ω4-def sv2' by simp
  subgoal by fact
  subgoal using ω4 unfolding ω4-def
  by (metis Opt.reach.Step ωω(2) fst-conv r(2) snd-conv)
  subgoal using ω4 unfolding ω4-def
  by (metis Nil-is-append-conv Van.reach-validFromS-reach last-snoc
not-Cons-self2 r(3))
  subgoal by fact subgoal by fact subgoal by fact subgoal by
fact
  subgoal using ωω by simp subgoal using ωω by simp subgoal
using ωω by simp ..
next

```

```

case False note trv2 = False
show ?thesis
apply(rule Van.llvalidFromS-selectlappend)
apply(rule exI[of - sv2]) apply(rule exI[of - trv2])
apply(rule exI[of - sv2']) apply(rule exI[of - w2'])
apply(rule exI[of - ltrv2']) apply(rule exI[of - n])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv2 .
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
    subgoal by fact
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Van.validFromS-def Van.validS-validTrans append-is-Nil-conv
list.sel(1) not-Cons-self2 trv2)
      subgoal apply(rule disjI1)
      apply(rule exI[of - trn'])
      apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule exI[of
- s1]) apply(rule exI[of - ltr1])
      apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])
      apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule
exI[of - statOO])
      apply(intro conjI)
      subgoal ..
      subgoal using trv2 unfolding ltrv2' trn' by auto
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
      subgoal by fact
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(2) snd-conv)
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
      subgoal by fact subgoal by fact subgoal by fact
      subgoal using  $\omega\omega$  by auto
      subgoal using  $\omega\omega$  by auto
      subgoal using  $\omega\omega$ 
      using llist-all-lappend-llist-of ltr1(3) by blast . .
qed
next
case False note current = current False
obtain w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO where
XX: XX s2 ltr2 tr2 s2' s2'' ltr2' and
 $\chi_4'$ :  $\chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2$ 
trv2 sv2'' statOO
and ltrv2: ltrv2 =

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```

      lappend (llist-of trv2) (lltrv2 (L, w1', w2', s1, ltr1, s2'', s2'' $ ltr2',
statA, sv1'', sv2'', statOO))
      using lltrv1-lltrv2-not-lnever-R[OF unw Δ r ltr2(1,2) isi3 ltr2(3)
current]
      unfolding ltrv2 by blast
      define ltrv2' where ltrv2': ltrv2' ≡ lltrv2 (L, w1', w2', s1, ltr1, s2'',
s2'' $ ltr2', statA, sv1'', sv2'', statOO)
      have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
      unfolding ltrv2 ltrv2' ..
      have trv2: trv2 ≠ [] using χ4' unfolding χ4'-def by auto

show ?thesis
apply(rule Van.llvalidFromS-selectlappend)
apply(rule exI[of - sv2]) apply(rule exI[of - trv2])
apply(rule exI[of - sv2''])
apply(rule exI[of - w2'])
apply(rule exI[of - ltrv2'])
apply(rule exI[of - n])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv2 .
  subgoal using χ4' unfolding χ4'-def
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
  subgoal using trv2 .
  subgoal using χ4' unfolding χ4'-def
  by (metis Simple-Transition-System.validFromS-def Van.validS-validTrans
append-is-Nil-conv list.sel(1) not-Cons-self2)
  subgoal apply(rule disjI1)
  apply(rule exI[of - L])
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
  apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
  apply(rule exI[of - statA]) apply(rule exI[of - sv1'']) apply(rule exI[of
- statOO])
  apply(intro conjI)
  subgoal .. subgoal unfolding ltrv2' ..
  subgoal using χ4' unfolding χ4'-def by simp
  subgoal by fact
  subgoal using χ4' unfolding χ4'-def
  by (metis Simple-Transition-System.reach-validFromS-reach χχ χχ-def

      append-is-Nil-conv last-snoc not-Cons-self2 r(2))
  subgoal using χ4' unfolding χ4'-def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using χ4' unfolding χ4'-def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)

```

```

      subgoal by fact
      subgoal by fact
      subgoal by fact
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of ltr2( $\beta$ ) by blast . .
    qed
  qed
  qed
  qed
  qed
}
thus ?thesis apply-apply(rule Van.lvalidFromS-imp-lvalidFromS)
using assms by blast
qed

```

```

lemma lcompletedFrom-lltrv1:
  assumes unw: unwindCond  $\Delta$ 
  and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
  shows Van.lcompletedFrom sv1 (lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO))
  proof-
    {fix ltrv1 assume ltrv1: ltrv1 = lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)
      and lfin: lfinite ltrv1
      hence list-of ltrv1  $\neq [] \wedge$  finalV (last (list-of ltrv1))
      using assms(2-) proof(induct length (list-of ltrv1) w1
        arbitrary: trn w2 ltrv1 s1 ltr1 s2 ltr2 statA sv1 sv2 statO
        rule: less2-induct')
      case (less w1 ltrv1 trn w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO)
        hence ltrv1: ltrv1 = lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
        statO)
          and lfin: lfinite ltrv1
          and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
          and r: reachO s1 reachO s2 reachV sv1 reachV sv2
          and ltr1: Opt.lvalidFromS s1 ltr1 lcompletedFromO s1 ltr1 lnever isIntO ltr1
          and ltr2: Opt.lvalidFromS s2 ltr2 lcompletedFromO s2 ltr2 lnever isIntO ltr2
          by auto
          have isi3:  $\neg$  isIntO s1 using ltr1
          by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
            llist.pred-inject(2))
          have isi4:  $\neg$  isIntO s2 using ltr2
          by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
            llist.pred-inject(2))
    }
  
```

```

show ?case proof(cases ltr1 = [] ∧ ltr2 = [])
  case True note ltr14 = True
    hence False using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by
auto
  thus ?thesis by auto
next
case False hence ltr14: ltr1 ≠ [] ∨ ltr2 ≠ [] by auto
show ?thesis proof(cases llength ltr1 ≤ Suc 0 ∨ llength ltr2 ≤ Suc 0)
  case True note ltr14 = ltr14 True
    hence ltrv1: list-of ltrv1 = [sv1] unfolding ltrv1 by simp
    have llength ltr1 = Suc 0 ∨ llength ltr2 = Suc 0
    using ltr14
    by (metis Opt.lcompletedFrom-def
      Suc-ile-eq i0-less lfinite-code(1) llength-eq-0 llist.exhaust
      ltr1(2) ltr2(2) nle-le not-lnull-conv zero-enat-def)
    hence ltr1 = [[s1]] ∨ ltr2 = [[s2]]
    using Opt.lcompletedFrom-singl ltr1(1) ltr1(2) ltr2(1) ltr2(2) by blast
    hence finalO s1 ∨ finalO s2
    using Opt.lcompletedFrom-LCons ltr1(2) ltr2(2) by blast
    hence finalV sv1
    using Δ r(1) r(2) r(3) r(4) unw unwindCond-def by auto
    thus ?thesis unfolding ltrv1 by auto
next
  case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0 by
auto
  show ?thesis
  proof(cases trn)
    case L note current = current L
    show ?thesis
    proof(cases lnever isSecO ltr1)
      case True note current = current True
        obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
        ωω: ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
        lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
        ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
        and trn' : trn' = (if trv1 = [] then L else R)
        and lltrv1: ltrv1 =
        lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
        using lltrv1-lltrv2-lnever-L[OF unw Δ r ltr1 isi4 current]
        unfolding ltrv1 by blast
        define ltrv1' where ltrv1': ltrv1' = lltrv1 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
        have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
        unfolding lltrv1 ltrv1' ..

        have trv1ne: trv1 ≠ [] ∨ w1' < w1 using ω3 unfolding ω3-def by
auto

```

```

have lfin': lfinite ltrv1'
using lfin trv1ne unfolding lltrv1 by simp
have len: length (list-of ltrv1') < length (list-of ltrv1) ∨
      length (list-of ltrv1') = length (list-of ltrv1) ∧ w1' < w1
using trv1ne lfin lfin' by (simp add: list-of-lappend lltrv1)

have 0: list-of ltrv1' ≠ [] ∧ finalV (last (list-of ltrv1'))
using len proof(elim disjE conjE)
  assume len: length (list-of ltrv1') < length (list-of ltrv1)
  show ?thesis
  apply(rule less(1)[OF - ltrv1'])
    subgoal by fact subgoal by fact
    subgoal using ω3 unfolding ω3-def by simp
    subgoal by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using ω3 unfolding ω3-def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using ω3 unfolding ω3-def
    by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
    subgoal by fact subgoal by fact subgoal by fact subgoal by fact
    subgoal by fact subgoal by fact .
  next
  assume len: length (list-of ltrv1') = length (list-of ltrv1) w1' < w1
  show ?thesis
  apply(rule less(2)[OF - - ltrv1'])
    subgoal by fact subgoal using len by simp subgoal by fact

    subgoal using ω3 unfolding ω3-def by simp
    subgoal by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using ω3 unfolding ω3-def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using ω3 unfolding ω3-def
    by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
    subgoal by fact subgoal by fact subgoal by fact subgoal by fact
    subgoal by fact subgoal by fact .
  qed
show ?thesis unfolding lltrv1 using 0
by (simp add: lfin' list-of-lappend)
next
case False note current = current False
obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
  XX: XX s1 ltr1 tr1 s1' s1'' ltr1' and
  X3': X3' Δ w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
and lltrv1: ltrv1 =
lappend (llist-of trv1) (lltrv1 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

```

```

using lltrv1-lltrv2-not-lnever-L[OF unw Δ r ltr1 isi4 current]
unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1' = lltrv1 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)

  have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding lltrv1 ltrv1' ..
  have trv1ne: trv1 ≠ [] using  $\chi^{\beta'}$  unfolding  $\chi^{\beta'-def}$  by auto
  have lfin': lfinite ltrv1'
  using lfin trv1ne unfolding lltrv1 by simp
  have len: length (list-of lltrv1') < length (list-of ltrv1)
  using trv1ne lfin lfin' by (simp add: list-of-lappend lltrv1)

  have 0: list-of ltrv1' ≠ [] ∧ finalV (last (list-of ltrv1'))
  apply(rule less(1)[OF - ltrv1'])
    subgoal by fact subgoal by fact
    subgoal using  $\chi^{\beta'}$  unfolding  $\chi^{\beta'-def}$  by simp
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$ 
      by (metis Simple-Transition-System.reach-validFromS-reach r(1)
snoc-eq-iff-butlast)
    subgoal by fact
    subgoal using  $\chi^{\beta'}$  unfolding  $\chi^{\beta'-def}$ 
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
r(3))
    subgoal using  $\chi^{\beta'}$  unfolding  $\chi^{\beta'-def}$ 
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by simp
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by simp
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$ 
      using llist-all-lappend-llist-of ltr1 by blast
    subgoal by fact subgoal by fact subgoal by fact .
  show ?thesis unfolding lltrv1 using 0
    by (simp add: lfin' list-of-lappend)
  qed
next
case R note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
  case True note current = current True
  obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
ωω: ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
ω4: ω4 Δ w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
  and trn': trn' = (if trv2 = [] then R else L)
  and ltrv1: ltrv1 =
lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,

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sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-R[OF un $\omega$   $\Delta$  r ltr2(1,2) isi3 ltr2(3) current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1' = lltrv1 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
  have lltrv1: ltrv1 = lappend (l $\omega$ list-of ltrv1) ltrv1'
  unfolding ltrv1 ltrv1' ..

  have trv1ne: ltrv1  $\neq$  []  $\vee$  w1' < w1 using  $\omega_4$  unfolding  $\omega_4$ -def by
auto
  have lfin': lfinite ltrv1'
  using lfin ltrv1ne unfolding lltrv1 by simp
  have len: length (l $\omega$ list-of ltrv1') < length (l $\omega$ list-of ltrv1)  $\vee$ 
length (l $\omega$ list-of ltrv1') = length (l $\omega$ list-of ltrv1)  $\wedge$  w1' < w1
  using trv1ne lfin lfin' by (simp add: l $\omega$ list-of-lappend lltrv1)

  have 0: l $\omega$ list-of ltrv1'  $\neq$  []  $\wedge$  finalV (last (l $\omega$ list-of ltrv1'))
  using len proof (elim disjE conjE)
  assume len: length (l $\omega$ list-of ltrv1') < length (l $\omega$ list-of ltrv1)
  show ?thesis
  apply (rule less(1)[OF - ltrv1'])
  subgoal by fact subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal by fact
  subgoal using r(2)  $\omega\omega$  by (metis Opt.reach.Step fst-conv snd-conv)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Van.reach-validFromS-reach append-is- $\omega$ Nil-conv last-snoc
not-Cons-self2 r(3))
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Van.reach-validFromS-reach append-is- $\omega$ Nil-conv last-snoc
not-Cons-self2 r(4))
  subgoal by fact subgoal by fact subgoal by fact subgoal by fact
  subgoal by fact subgoal by fact .
next
  assume len: length (l $\omega$ list-of ltrv1') = length (l $\omega$ list-of ltrv1) w1' < w1
  show ?thesis
  apply (rule less(2)[OF - - ltrv1'])
  subgoal by fact subgoal using len by simp subgoal by fact

  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal by fact
  subgoal by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(2) snd-conv)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Van.reach-validFromS-reach append-is- $\omega$ Nil-conv last-snoc
not-Cons-self2 r(4))
  subgoal by fact subgoal by fact subgoal by fact subgoal by fact
  subgoal by fact subgoal by fact .

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qed
show ?thesis unfolding lltrv1 using 0
by (simp add: lfin' list-of-lappend)
next
case False note current = current False
obtain w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO where
  XX: XX s2 ltr2 tr2 s2' s2'' ltr2' and
   $\chi_4'$ :  $\chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2$ 
trv2 sv2'' statOO
and ltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (L, w1', w2', s1, ltr1, s2'', s2'' $ ltr2',
statA, sv1'', sv2'', statOO))
  using lltrv1-lltrv2-not-lnever-R[OF unW  $\Delta$  r ltr2(1,2) isi3 ltr2(3)
current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1' = lltrv1 (L, w1', w2', s1, ltr1, s2'',
s2'' $ ltr2', statA, sv1'', sv2'', statOO)
  have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..

  have trv1ne: trv1  $\neq$  []  $\vee$  w1' < w1 using  $\chi_4'$  unfolding  $\chi_4'$ -def by
auto

  have lfin': lfinite ltrv1'
  using lfin trv1ne unfolding lltrv1 by simp
  have len: length (list-of ltrv1') < length (list-of ltrv1)  $\vee$ 
    length (list-of ltrv1') = length (list-of ltrv1)  $\wedge$  w1' < w1
  using trv1ne lfin lfin' by (simp add: list-of-lappend lltrv1)

  have 0: list-of ltrv1'  $\neq$  []  $\wedge$  finalV (last (list-of ltrv1'))
  using len proof (elim disjE conjE)
  assume len: length (list-of ltrv1') < length (list-of ltrv1)
  show ?thesis
  apply (rule less(1)[OF - ltrv1'])
  subgoal by fact subgoal by fact
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
  subgoal by fact
  subgoal using r(2) XX unfolding XX-def
  by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
  by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(3))
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
  by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
  subgoal by fact subgoal by fact subgoal by fact
  subgoal using XX unfolding XX-def by auto
  subgoal using XX unfolding XX-def by auto
  subgoal using XX unfolding XX-def

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    using llist-all-lappend-llist-of-ltr2(3) by blast .
  next
  assume len: length (list-of ltrv1') = length (list-of ltrv1) w1' < w1
  show ?thesis
  apply(rule less(2)[OF - - ltrv1'])
    subgoal by fact subgoal using len by simp subgoal by fact

    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
    subgoal by fact
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(2))
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
r(4))

    subgoal by fact subgoal by fact subgoal by fact
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of-ltr2(3) by blast .
  qed
  show ?thesis unfolding lltrv1 using 0
  by (simp add: lfn' list-of-lappend)
  qed
  qed
  qed
  qed
  qed
  }
  thus ?thesis unfolding Van.lcompletedFrom-def by auto
qed

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lemma lcompletedFrom-lltrv2:
  assumes unw: unwindCond  $\Delta$ 
  and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
  shows Van.lcompletedFrom sv2 (lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  proof -
  {fix ltrv2 assume ltrv2: ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and lfn: lfinite ltrv2
  hence list-of ltrv2  $\neq [] \wedge$  finalV (last (list-of ltrv2))
  using assms(2-) proof(induct length (list-of ltrv2) w2
  arbitrary: ltrv2 trn w1 s1 ltr1 s2 ltr2 statA sv1 sv2 statO
  rule: less2-induct')
  case (less w2 ltrv2 trn w1 s1 ltr1 s2 ltr2 statA sv1 sv2 statO)

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hence ltrv2: ltrv2 = lltrv2 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
and lfin: lfinite ltrv2
and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr1: Opt.lvalidFromS s1 ltr1 lcompletedFromO s1 ltr1 lnever isIntO ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 lcompletedFromO s2 ltr2 lnever isIntO ltr2
by auto
have isi3: ¬ isIntO s1 using ltr1
by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
llist.pred-inject(2))
have isi4: ¬ isIntO s2 using ltr2
by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
llist.pred-inject(2))

show ?case proof(cases ltr1 = [[]] ∧ ltr2 = [[]])
case True note ltr14 = True
hence False using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by
auto
thus ?thesis by auto
next
case False hence ltr14: ltr1 ≠ [[]] ∨ ltr2 ≠ [[]] by auto
show ?thesis proof(cases llength ltr1 ≤ Suc 0 ∨ llength ltr2 ≤ Suc 0)
case True note ltr14 = ltr14 True
hence ltrv2: list-of ltrv2 = [sv2] unfolding ltrv2 by simp
have llength ltr1 = Suc 0 ∨ llength ltr2 = Suc 0
using ltr14
by (metis Opt.lcompletedFrom-def
Suc-ile-eq i0-less lfinite-code(1) llength-eq-0 llist.exhaust
ltr1(2) ltr2(2) nle-le not-lnull-conv zero-enat-def)
hence ltr1 = [[s1]] ∨ ltr2 = [[s2]]
using Opt.lcompletedFrom-singl ltr1(1) ltr1(2) ltr2(1) ltr2(2) by blast
hence finalO s1 ∨ finalO s2
using Opt.lcompletedFrom-LCons ltr1(2) ltr2(2) by blast
hence finalV sv2
using Δ r(1) r(2) r(3) r(4) unw unwindCond-def by auto
thus ?thesis unfolding ltrv2 by auto
next
case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0 by
auto
show ?thesis
proof(cases trn)
case L note current = current L
show ?thesis
proof(cases lnever isSecO ltr1)
case True note current = current True
obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
ωω: ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
lcompletedFromO s1' ltr1' lnever isIntO ltr1' and

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statOO      ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
and trn' : trn' = (if trv1 = [] then L else R)
and lltrv2: ltrv2 =
  lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
using lltrv1-lltrv2-lnever-L[OF unω Δ r ltr1 isi4 current]
unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2' = lltrv2 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
  have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding lltrv2 ltrv2' ..

  have trv2ne: trv2 ≠ [] ∨ w2' < w2 using ω3 unfolding ω3-def by
auto
  have lfin': lfinite ltrv2'
  using lfin trv2ne unfolding lltrv2 by simp
  have len: length (list-of ltrv2') < length (list-of ltrv2) ∨
    length (list-of ltrv2') = length (list-of ltrv2) ∧ w2' < w2
  using trv2ne lfin lfin' by (simp add: list-of-lappend lltrv2)

  have 0: list-of ltrv2' ≠ [] ∧ finalV (last (list-of ltrv2'))
  using len proof (elim disjE conjE)
  assume len: length (list-of ltrv2') < length (list-of ltrv2)
  show ?thesis
  apply (rule less(1)[OF - ltrv2'])
  subgoal by fact subgoal by fact
  subgoal using ω3 unfolding ω3-def by simp
  subgoal by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
  subgoal by fact
  subgoal using ω3 unfolding ω3-def
  by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
  subgoal using ω3 unfolding ω3-def
  by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
  subgoal by fact subgoal by fact subgoal by fact subgoal by fact
  subgoal by fact subgoal by fact .
next
assume len: length (list-of ltrv2') = length (list-of ltrv2) w2' < w2
show ?thesis
apply (rule less(2)[OF - - ltrv2'])
  subgoal by fact subgoal using len by simp subgoal by fact

  subgoal using ω3 unfolding ω3-def by simp
  subgoal by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
  subgoal by fact
  subgoal using ω3 unfolding ω3-def
  by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
  subgoal using ω3 unfolding ω3-def

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      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
      subgoal by fact subgoal by fact subgoal by fact subgoal by fact
      subgoal by fact subgoal by fact .
    qed
    show ?thesis unfolding lltrv2 using 0
    by (simp add: lfin' list-of-lappend)
  next
    case False note current = current False
    obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
     $\chi\chi$ :  $\chi\chi$  s1 ltr1 tr1 s1' s1'' ltr1' and
     $\chi^3$ ':  $\chi^3$ '  $\Delta$  w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
    and lltrv2: ltrv2 =
lappend (llist-of trv2) (lltrv2 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

    using lltrv1-lltrv2-not-lnever-L[OF unW  $\Delta$  r ltr1 isi4 current]
    unfolding ltrv2 by blast
    define ltrv2' where ltrv2': ltrv2' = lltrv2 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
    have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
    unfolding lltrv2 ltrv2' ..

    have trv2ne: trv2  $\neq$  []  $\vee$  w2' < w2 using  $\chi^3$ ' unfolding  $\chi^3$ '-def by
auto
    have lfin': lfinite ltrv2'
    using lfin trv2ne unfolding lltrv2 by simp
    have len: length (list-of ltrv2') < length (list-of ltrv2)  $\vee$ 
length (list-of ltrv2') = length (list-of ltrv2)  $\wedge$  w2' < w2
    using trv2ne lfin lfin' by (simp add: list-of-lappend lltrv2)

    have 0: list-of ltrv2'  $\neq$  []  $\wedge$  finalV (last (list-of ltrv2'))
    using len proof (elim disjE conjE)
      assume len: length (list-of ltrv2') < length (list-of ltrv2)
      show ?thesis
      apply (rule less(1)[OF - ltrv2'])
      subgoal by fact subgoal by fact
      subgoal using  $\chi^3$ ' unfolding  $\chi^3$ '-def by simp
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(1)
snoc-eq-iff-butlast)
      subgoal by fact
      subgoal using  $\chi^3$ ' unfolding  $\chi^3$ '-def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
r(3))
      subgoal using  $\chi^3$ ' unfolding  $\chi^3$ '-def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by simp

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    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by simp
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
    using llist-all-lappend-llist-of ltr1 by blast
    subgoal by fact subgoal by fact subgoal by fact .
next
assume len: length (list-of ltrv2') = length (list-of ltrv2) w2' < w2
show ?thesis
apply(rule less(2)[OF - - ltrv2'])
    subgoal by fact subgoal using len by simp subgoal by fact

    subgoal using  $\chi\beta'$  unfolding  $\chi\beta'$ -def by simp
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(1))
    subgoal by fact
    subgoal using  $\chi\beta'$  unfolding  $\chi\beta'$ -def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using  $\chi\beta'$  unfolding  $\chi\beta'$ -def
    by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
    using llist-all-lappend-llist-of ltr1(3) by blast
    subgoal by fact subgoal by fact subgoal by fact .
qed
show ?thesis unfolding lltrv2 using 0
by (simp add: lfn' list-of-lappend)
qed
next
case R note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
case True note current = current True
obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
 $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
 $\omega_4$ :  $\omega_4 \Delta w1 w2 w1' w2' s1 s2 s2'$  statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
and trn': trn' = (if trv2 = [] then R else L)
and ltrv2: ltrv2 =
lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
using lltrv1-lltrv2-lnever-R[OF unW  $\Delta r$  ltr2(1,2) isi3 ltr2(3) current]
unfolding ltrv2 by blast
define ltrv2' where ltrv2': ltrv2' = lltrv2 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
unfolding ltrv2 ltrv2' ..

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      have trv2ne: trv2 ≠ [] ∨ w2' < w2 using ω4 unfolding ω4-def by
auto
      have lfin': lfinite ltrv2'
      using lfin trv2ne unfolding lltrv2 by simp
      have len: length (list-of ltrv2') < length (list-of ltrv2) ∨
        length (list-of ltrv2') = length (list-of ltrv2) ∧ w2' < w2
      using trv2ne lfin lfin' by (simp add: list-of-lappend lltrv2)

      have 0: list-of ltrv2' ≠ [] ∧ finalV (last (list-of ltrv2'))
      using len proof (elim disjE conjE)
        assume len: length (list-of ltrv2') < length (list-of ltrv2)
        show ?thesis
        apply (rule less(1)[OF - ltrv2'])
          subgoal by fact subgoal by fact
          subgoal using ω4 unfolding ω4-def by simp
          subgoal by fact
          subgoal using r(2) ωω by (metis Opt.reach.Step fst-conv snd-conv)
          subgoal using ω4 unfolding ω4-def
            by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(3))
          subgoal using ω4 unfolding ω4-def
            by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
          subgoal by fact subgoal by fact subgoal by fact subgoal by fact
          subgoal by fact subgoal by fact .
        next
        assume len: length (list-of ltrv2') = length (list-of ltrv2) w2' < w2
        show ?thesis
        apply (rule less(2)[OF - - ltrv2'])
          subgoal by fact subgoal using len by simp subgoal by fact

          subgoal using ω4 unfolding ω4-def by simp
          subgoal by fact
          subgoal by (metis Opt.reach.Step ωω(2) fst-conv r(2) snd-conv)
          subgoal using ω4 unfolding ω4-def
            by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
          subgoal using ω4 unfolding ω4-def
            by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
          subgoal by fact subgoal by fact subgoal by fact subgoal by fact
          subgoal by fact subgoal by fact .
        qed
      show ?thesis unfolding lltrv2 using 0
      by (simp add: lfin' list-of-lappend)
    next
    case False note current = current False
    obtain w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO where
    XX: XX s2 ltr2 tr2 s2' s2'' ltr2' and

```

```

       $\chi_4'$ :  $\chi_4' \Delta w_1 w_2 w_1' w_2' s_1 s_2 tr_2 s_2' s_2'' statA sv_1 trv_1 sv_1'' sv_2$ 
\Delta r ltr2(1,2) isi3 ltr2(3)
current]
      unfolding ltrv2 by blast
      define ltrv2' where ltrv2': ltrv2' = lltrv2 (L, w1', w2', s1, ltr1, s2'',
s2'' $ ltr2', statA, sv1'', sv2'', statOO)
      have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
      unfolding ltrv2 ltrv2' ..

      have trv2ne: trv2  $\neq$  [] using  $\chi_4'$  unfolding  $\chi_4'$ -def by auto
      have lfin': lfinite ltrv2'
      using lfin trv2ne unfolding lltrv2 by simp
      have len: length (list-of ltrv2') < length (list-of ltrv2)
      using trv2ne lfin lfin' by (simp add: list-of-lappend lltrv2)

      have 0: list-of ltrv2'  $\neq$  []  $\wedge$  finalV (last (list-of ltrv2'))
      apply(rule less(1)[OF - ltrv2'])
      subgoal by fact subgoal by fact
      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
      subgoal by fact
      subgoal using r(2)  $\chi\chi$  unfolding  $\chi\chi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(3))
      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
r(4))

      subgoal by fact subgoal by fact subgoal by fact
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of ltr2(3) by blast .
      show ?thesis unfolding lltrv2 using 0
      by (simp add: lfin' list-of-lappend)
      qed
    qed
  qed
  qed
}
}
thus ?thesis unfolding Van.lcompletedFrom-def by auto
qed

```


lemma *lS-lltrv1-ltr1*:
assumes *unw*: *unwindCond* Δ
and Δ : $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*
and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1* *lnever isIntO* *ltr1*
and *ltr2*: *Opt.lvalidFromS* *s2* *ltr2* *Opt.lcompletedFrom* *s2* *ltr2* *lnever isIntO* *ltr2*
shows *Van.lS* (*lltrv1* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*)) = *Opt.lS* *ltr1*

proof–
have *cltrv1*: *Van.lcompletedFrom* *sv1* (*lltrv1* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*))
using *lcompletedFrom-lltrv1* [*OF assms*] .
{fix *trn* *nL* *nR* *ltrv1* *ltr1*
assume $\exists w1\ w2\ s1\ s2\ ltr2\ statA\ sv1\ sv2\ statO.$
 $nL = w1 \wedge nR = w2 \wedge$
 $ltrv1 = lltrv1\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$
 $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$
 $reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$
 $Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge lnever\ isIntO\ ltr1 \wedge$
 $Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge lnever\ isIntO\ ltr2$
hence *TwoFuncPred.sameFM1 isSecV isSecO getSecV getSecO* *trn* *nL* *nR* *ltrv1* *ltr1*
proof(*coinduct* rule: *TwoFuncPred.sameFM1.coinduct*[of $\lambda trn\ nL\ nR\ ltrv1\ ltr1.$

$\exists w1\ w2\ s1\ s2\ ltr2\ statA\ sv1\ sv2\ statO.$
 $nL = w1 \wedge nR = w2 \wedge$
 $ltrv1 = lltrv1\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$
 $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$
 $reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$
 $Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge lnever\ isIntO\ ltr1 \wedge$
 $Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge lnever\ isIntO\ ltr2,$
where *pred* = *isSecV* **and** *pred'* = *isSecO* **and** *func* = *getSecV* **and** *func'*
= *getSecO*])
case (*?* *trn* *nL* *nR* *ltrv1* *ltr1*)
then obtain *w1* *w2* *sv1* *s1* *s2* *ltr2* *statA* *sv2* *statO*
where *nL*: *nL* = *w1* **and** *nR*: *nR* = *w2*
and *ltrv1*: *ltrv1* = *lltrv1* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*)
and Δ : $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*
and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1* *lnever isIntO* *ltr1*
and *ltr2*: *Opt.lvalidFromS* *s2* *ltr2* *Opt.lcompletedFrom* *s2* *ltr2* *lnever isIntO* *ltr2*
by *auto*
have *isi3*: $\neg isIntO\ s1$ **using** *ltr1*
by (*metis* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-def* *lfinite-LNil* *llist.exhaust-sel* *llist.pred-inject*(*?*))
have *isi4*: $\neg isIntO\ s2$ **using** *ltr2*
by (*metis* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-def* *lfinite-LNil* *llist.exhaust-sel* *llist.pred-inject*(*?*))

```

show ?case proof(cases ltr1 = [] ∧ ltr2 = [])
  case True note ltr14 = True
  hence ltrv1: ltrv1 = [] unfolding ltrv1 by simp
show ?thesis using ltr14 unfolding ltrv1 apply-apply(rule TwoFuncPred.sameFM1-selectLNil)
by auto
next
case False hence ltr14: ltr1 ≠ [] ∨ ltr2 ≠ [] by auto
show ?thesis proof(cases llength ltr1 ≤ Suc 0 ∨ llength ltr2 ≤ Suc 0)
  case True note ltr14 = ltr14 True
  hence ltrv1: ltrv1 = [[sv1]] unfolding ltrv1 by simp
  have llength ltr1 = Suc 0 ∨ llength ltr2 = Suc 0
  by (metis Opt.lcompletedFrom-def Suc-ile-eq True
    lfinite-LNil llength-LNil llist-eq-cong ltr1(2)
    ltr2(2) nle-le order-le-imp-less-or-eq zero-enat-def zero-order(3))
  hence finalO s1 ∨ finalO s2
  using Opt.lcompletedFrom-singl ltr1(1) ltr1(2) ltr2(1) ltr2(2) by blast
  hence fs1: finalO s1
  using Δ r(1) r(2) r(3) r(4) unω unwindCond-def by auto
  hence ltr1: ltr1 = [[s1]]
  by (metis Opt.final-def Opt.lcompletedFrom-def
    Opt.lvalidFromS-Cons-iff lfinite-code(1) llist.exhaust ltr1(1) ltr1(2))
  have fsv1: finalV sv1
  using Δ fs1 r(1) r(2) r(3) r(4) unω unwindCond-final by blast
  have isv13: ¬ isSecV sv1 ∧ ¬ isSecO s1
  using fsv1 fs1 Opt.final-not-isSec Van.final-not-isSec by blast
show ?thesis unfolding ltrv1 ltr1 apply(rule TwoFuncPred.sameFM1-selectSingl)

  using isv13 by auto
next
case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0
by auto
show ?thesis proof(cases trn)
  case L note trn = L[simp] note current = current L
  show ?thesis
  proof(cases lnever isSecO ltr1)
    case True note current = current True
    obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
      ωω: ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
      lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
      ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
    and trn': trn' = (if trv1 = [] then L else R)
    and lltrv1: lltrv1 =
      lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
    using lltrv1-lltrv2-lnever-L[OF unω Δ r ltr1 isi4 current]
    unfolding ltrv1 by blast
    define ltrv1' where ltrv1': ltrv1' ≡ lltrv1 (trn', w1', w2', s1', ltr1',

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s2, ltr2, statA, sv1', sv2', statOO)
  have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..
  have nis1: ¬ isSecO s1 using True ωω(1) by force
  show ?thesis
  proof(cases trv1 = [])
    case True note trv1 = True
    hence w1' < w1 using ω3 unfolding ω3-def by auto
    have [simp]: trn' = trn by (simp add: trv1 trn')
    show ?thesis
    apply(rule TwoFuncPred.sameFM1-selectDelayL)
    apply(rule exI[of - w1']) apply(rule exI[of - w1])
    apply(rule exI[of - trv1]) apply(rule exI[of - [s1]])
    apply(rule exI[of - w2'])
    apply(rule exI[of - ltrv1']) apply(rule exI[of - ltr1'])
    apply(rule exI[of - w2])
    apply(intro conjI)
    subgoal by fact
    subgoal unfolding nL .. subgoal unfolding nR ..
    subgoal unfolding ltrv1 trv1 by simp
    subgoal unfolding ωω(1) by simp
    subgoal by fact subgoal unfolding trv1 using ω3-def nis1 by
simp
  subgoal apply(rule disjI1)
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1']) apply(rule exI[of - s2])
    apply(rule exI[of - ltr2]) apply(rule exI[of - statA])
    apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
    apply(rule exI[of - statOO])
    apply(intro conjI)
    subgoal .. subgoal ..
    subgoal unfolding ltrv1' by simp
    subgoal using ω3 unfolding ω3-def by simp
    subgoal using ω3 unfolding ω3-def
    by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using ω3 unfolding ω3-def
    by (metis Simple-Transition-System.reach-validFromS-reach r(3))
  snoc-eq-iff-butlast)
  subgoal using ω3 unfolding ω3-def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4))
  snoc-eq-iff-butlast)
  subgoal using ωω by auto
  subgoal using ωω by auto
  subgoal using ωω
  using llist-all-lappend-llist-of ltr1(3) by blast
  subgoal using ωω using ltr2(1) by fastforce
  subgoal by fact
  subgoal by fact . .

```

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next
  case False note trv1 = False
  show ?thesis
  apply(rule TwoFuncPred.sameFM1-selectlappend)
  apply(rule exI[of - trv1]) apply(rule exI[of - [s1]])
  apply(rule exI[of - trn']) apply(rule exI[of - w1'])
  apply(rule exI[of - w2'])
  apply(rule exI[of - ltrv1']) apply(rule exI[of - ltr1'])
  apply(rule exI[of - trn])
  apply(rule exI[of - w1])
  apply(rule exI[of - w2])
  apply(intro conjI)
  subgoal ..
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrv1 .
  subgoal unfolding  $\omega\omega(1)$  by simp
  subgoal by fact
  subgoal using  $\omega\exists$  unfolding  $\omega\exists$ -def by simp
  subgoal using ltr1(3) \omega\exists unfolding  $\omega\exists$ -def
    by (metis Opt.S.map-filter Opt.S.simps(4) Van.S.map-filter
Van.S.eq-Nil-iff(2) append-Nil
butlast-snoc filter.simps(2) nis1)
  subgoal apply(rule disjI1)
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1']) apply(rule exI[of - s2])
    apply(rule exI[of - ltr2]) apply(rule exI[of - statA])
    apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
    apply(rule exI[of - statOO])
    apply(intro conjI)
    subgoal .. subgoal ..
    subgoal unfolding ltrv1' ..
    subgoal using  $\omega\exists$  unfolding  $\omega\exists$ -def by simp
    subgoal using  $\omega\exists$  unfolding  $\omega\exists$ -def
      by (metis Opt.reach.Step \omega\omega(2) fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using  $\omega\exists$  unfolding  $\omega\exists$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
    subgoal using  $\omega\exists$  unfolding  $\omega\exists$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$ 
      using l1ist-all-lappend-l1ist-of ltr1(3) by blast
    subgoal using  $\omega\omega$  using ltr2(1) by fastforce
    subgoal by fact
    subgoal by fact . .
  qed

```

```

next
  case False note current = current False
  obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
  χχ: χχ s1 ltr1 tr1 s1' s1'' ltr1' and
  χ3': χ3' Δ w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
  and lltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

  using lltrv1-lltrv2-not-lnever-L[OF unW Δ r ltr1 isi4 current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1' ≡ lltrv1 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
  have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding lltrv1 ltrv1' ..

show ?thesis apply(rule TwoFuncPred.sameFM1-selectlappend)
apply(rule exI[of - trv1]) apply(rule exI[of - tr1 ## s1'])
apply(rule exI[of - R])
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - ltrv1']) apply(rule exI[of - s1'' $ ltr1'])
apply(rule exI[of - trn])
apply(rule exI[of - w1]) apply(rule exI[of - w2])
apply(intro conjI)
  subgoal .. subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrv1 .
  subgoal using χχ unfolding χχ-def by simp
  subgoal using χ3' unfolding χ3'-def by simp
  subgoal by simp
  subgoal using χ3' unfolding χ3'-def
  by (simp add: Opt.S.map-filter Van.S.map-filter)
  subgoal apply(rule disjI1)
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1''])
  apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
  apply(rule exI[of - statA]) apply(rule exI[of - sv1'']) apply(rule exI[of
- sv2''])

  apply(rule exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1' ..
  subgoal using χ3' unfolding χ3'-def by simp
  subgoal using χ3' unfolding χ3'-def
  by (metis Simple-Transition-System.reach-validFromS-reach χχ χχ-def

  append-is-Nil-conv last-snoc not-Cons-self2 r(1))
  subgoal by fact
  subgoal using χ3' unfolding χ3'-def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3))

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snoc-eq-iff-butlast)
  subgoal using  $\chi\beta'$  unfolding  $\chi\beta'$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach r(4))
snoc-eq-iff-butlast)
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
    using llist-all-lappend-llist-of ltr1(3) by blast
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def using ltr2(1) by fastforce
  subgoal by fact
  subgoal by fact . .
qed
next
case R note trn = R[simp] note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
  case True note current = current True
  obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
   $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
  lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
   $\omega_4$ :  $\omega_4 \Delta w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'$ 
statOO
  and trn': trn' = (if trv2 = [] then R else L)
  and ltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-R[OF unW  $\Delta$  r ltr2(1,2) isi3 ltr2(3) current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
  have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..
  have nev1: never isSecV trv1 using  $\omega_4$  unfolding  $\omega_4$ -def by auto
  show ?thesis
proof(cases trv2 = [])
  case True note trv2 = True
  have [simp]: trn' = trn using R trv2 trn' by auto
  have w2' < w2 using  $\omega_4$  trv2 unfolding  $\omega_4$ -def by auto
  show ?thesis
  apply(rule TwoFuncPred.sameFM1-selectDelayR)
  apply(rule exI[of - w2']) apply(rule exI[of - nR])
  apply(rule exI[of - trv1]) apply(rule exI[of - []])
  apply(rule exI[of - w1'])
  apply(rule exI[of - ltrv1']) apply(rule exI[of - ltr1])
  apply(rule exI[of - nL])
  apply(intro conjI)
  subgoal by simp subgoal .. subgoal ..
  subgoal by fact subgoal by simp
  subgoal unfolding nR by fact

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subgoal using nev1 by (simp add: never-Nil-filter)
subgoal apply(rule disjI1)
  apply(rule exI[of - w1 ^]) apply(rule exI[of - w2 ^])
  apply(rule exI[of - s1]) apply(rule exI[of - s2 ^])
  apply(rule exI[of - ltr2 ^]) apply(rule exI[of - statA])
  apply(rule exI[of - sv1 ^]) apply(rule exI[of - sv2 ^])
  apply(rule exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1' by simp
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(2) snd-conv)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Van.reach-validFromS-reach r(4) snoc-eq-iff-butlast)
  subgoal by fact subgoal by fact subgoal by fact
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto . .
next
case False note trv2 = False
have [simp]: trn' = L using R trv2 trn' by auto
show ?thesis
apply(rule TwoFuncPred.sameFM1-selectRL)
apply(rule exI[of - trv1]) apply(rule exI[of - []])
apply(rule exI[of - w1 ^]) apply(rule exI[of - w2 ^])
apply(rule exI[of - ltrv1 ^]) apply(rule exI[of - ltr1])
apply(rule exI[of - w1]) apply(rule exI[of - w2])
apply(intro conjI)
  subgoal by fact
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal unfolding ltrv1 ..
  subgoal unfolding  $\omega\omega(1)$  by simp
subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by (simp add: never-Nil-filter)
  subgoal apply(rule disjI1)
  apply(rule exI[of - w1 ^]) apply(rule exI[of - w2 ^])
  apply(rule exI[of - s1]) apply(rule exI[of - s2 ^])
  apply(rule exI[of - ltr2 ^]) apply(rule exI[of - statA])
  apply(rule exI[of - sv1 ^]) apply(rule exI[of - sv2 ^])
  apply(rule exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1' by simp
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def

```

```

by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(2) snd-conv)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Van.reach-validFromS-reach r(4) snoc-eq-iff-butlast)
  subgoal by fact subgoal by fact subgoal by fact
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto . .

qed
next
case False note current = current False
obtain  $w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO$  where
 $\chi\chi: \chi\chi s2 ltr2 tr2 s2' s2'' ltr2'$  and
 $\chi_4': \chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2$ 
 $trv2 sv2'' statOO$ 
and  $ltrv1: ltrv1 =$ 
  lappend (llist-of  $trv1$ ) (lltrv1 (L,  $w1', w2', s1, ltr1, s2'', s2'' \$ ltr2',$ 
 $statA, sv1'', sv2'', statOO$ ))
  using lltrv1-lltrv2-not-lnever-R[OF  $unw \Delta r ltr2(1,2) isi3 ltr2(3)$ 
current]
  unfolding  $ltrv1$  by blast
  define  $ltrv1'$  where  $ltrv1': ltrv1' \equiv lltrv1 (L, w1', w2', s1, ltr1, s2'',$ 
 $s2'' \$ ltr2', statA, sv1'', sv2'', statOO)$ 
  have  $ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'$ 
  unfolding  $ltrv1 ltrv1' ..$ 

show ?thesis
apply(rule TwoFuncPred.sameFM1-selectRL)
apply(rule exI[of -  $trv1$ ]) apply(rule exI[of - []])
apply(rule exI[of -  $w1 \uparrow$ ]) apply(rule exI[of -  $w2 \uparrow$ ])
apply(rule exI[of -  $ltrv1 \uparrow$ ]) apply(rule exI[of -  $ltr1$ ])
apply(rule exI[of -  $w1$ ]) apply(rule exI[of -  $w2$ ])
apply(intro conjI)
  subgoal by fact
  subgoal unfolding  $nL ..$  subgoal unfolding  $nR ..$ 
  subgoal unfolding  $ltrv1 ..$  subgoal by simp
subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by (simp add: never-Nil-filter)
  subgoal apply(rule disjI1)
  apply(rule exI[of -  $w1 \uparrow$ ]) apply(rule exI[of -  $w2 \uparrow$ ])
  apply(rule exI[of -  $s1$ ]) apply(rule exI[of -  $s2' \uparrow$ ])
  apply(rule exI[of -  $s2'' \$ ltr2 \uparrow$ ]) apply(rule exI[of -  $statA$ ])
  apply(rule exI[of -  $sv1' \uparrow$ ]) apply(rule exI[of -  $sv2' \uparrow$ ])
  apply(rule exI[of -  $statOO$ ])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding  $ltrv1'$  by simp
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
  subgoal by fact

```



```

      subgoal using r(2)  $\chi\chi$  unfolding  $\chi\chi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
      subgoal using r(3)  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
      subgoal using r(4)  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
      subgoal by fact subgoal by fact subgoal by fact
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of ltr2(3) by blast . .
    qed
  qed
  qed
  qed
  qed
}
thus ?thesis unfolding Van.lS[OF cltrv1] Opt.lS[OF ltr1(2)]
apply- apply(rule TwoFuncPred.sameFM1-lmap-lfilter)
using assms by blast
qed

```

lemma *lS-lltrv2-ltr2*:

assumes *unw*: *unwindCond* Δ

and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$

and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*

and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1* *lnever isIntO* *ltr1*

and *ltr2*: *Opt.lvalidFromS* *s2* *ltr2* *Opt.lcompletedFrom* *s2* *ltr2* *lnever isIntO* *ltr2*

shows *Van.lS* (*lltrv2* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*)) = *Opt.lS* *ltr2*

proof–

have *cltrv2*: *Van.lcompletedFrom* *sv2* (*lltrv2* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*))

using *lcompletedFrom-lltrv2*[*OF* *assms*] .

{fix *trn* *nL* *nR* *ltrv2* *ltr2*

assume $\exists w1 w2 s1 s2 ltr1 statA sv1 sv2 statO$.

$nL = w1 \wedge nR = w2 \wedge$

$ltrv2 = lltrv2 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$

$\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$

$reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$

$Opt.lvalidFromS s1 ltr1 \wedge Opt.lcompletedFrom s1 ltr1 \wedge lnever isIntO ltr1 \wedge$

$Opt.lvalidFromS s2 ltr2 \wedge Opt.lcompletedFrom s2 ltr2 \wedge lnever isIntO ltr2$

hence *TwoFuncPred.sameFM2* *isSecV* *isSecO* *getSecV* *getSecO* *trn* *nL* *nR* *ltrv2* *ltr2*

proof(*coinduct* rule: *TwoFuncPred.sameFM2.coinduct*[of $\lambda trn nL nR ltrv2 ltr2$.

$\exists w1 w2 s1 s2 ltr1 statA sv1 sv2 statO$.

$nL = w1 \wedge nR = w2 \wedge$

```

    ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ∧
    Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
    reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
    Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧ lnever isIntO ltr1 ∧

    Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧ lnever isIntO ltr2,
    where pred = isSecV and pred' = isSecO and func = getSecV and func'
= getSecO])
  case (2 trn nL nR ltrv2 ltr2)
  then obtain w1 w2 sv1 s1 s2 ltr1 statA sv2 statO
  where nL: nL = w1 and nR: nR = w2
  and ltrv2: ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
  by auto
  have isi3: ¬ isIntO s1 using ltr1
  by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
llist.pred-inject(2))
  have isi4: ¬ isIntO s2 using ltr2
  by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
llist.pred-inject(2))

  show ?case proof(cases ltr1 = [] ∧ ltr2 = [])
  case True note ltr14 = True
  hence ltrv2: ltrv2 = [] unfolding ltrv2 by simp
  show ?thesis using ltr14 unfolding ltrv2 apply-apply(rule TwoFuncPred.sameFM2-selectLNil)
by auto
next
  case False hence ltr14: ltr1 ≠ [] ∨ ltr2 ≠ [] by auto
  show ?thesis proof(cases llength ltr1 ≤ Suc 0 ∨ llength ltr2 ≤ Suc 0)
  case True note ltr14 = ltr14 True
  hence ltrv2: ltrv2 = [[sv2]] unfolding ltrv2 by simp
  have llength ltr1 = Suc 0 ∨ llength ltr2 = Suc 0
  by (metis Opt.lcompletedFrom-def Suc-ile-eq True
lfinite-LNil llength-LNil llist-eq-cong ltr1(2)
ltr2(2) nle-le order-le-imp-less-or-eq zero-enat-def zero-order(3))
  hence finalO s1 ∨ finalO s2
  using Opt.lcompletedFrom-singl ltr1(1) ltr1(2) ltr2(1) ltr2(2) by blast
  hence fs2: finalO s2
  using Δ r(1) r(2) r(3) r(4) unW unwindCond-def by auto
  hence ltr2: ltr2 = [[s2]]
  by (metis Opt.final-def Opt.lcompletedFrom-def
Opt.lvalidFromS-Cons-iff lfinite-code(1) llist.exhaust ltr2(1) ltr2(2))
  have fsv2: finalV sv2
  using Δ fs2 r(1) r(2) r(3) r(4) unW unwindCond-final by blast
  have isv24: ¬ isSecV sv2 ∧ ¬ isSecO s2
  using fsv2 fs2 Opt.final-not-isSec Van.final-not-isSec by blast

```

```

show ?thesis unfolding ltrv2 ltr2 apply(rule TwoFuncPred.sameFM2-selectSingl)

  using isv24 by auto
next
  case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0
  by auto
  show ?thesis proof(cases trn)
    case L note trn = L[simp] note current = current L
    show ?thesis
    proof(cases lnever isSecO ltr1)
      case True note current = current True
      obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
        ωω: ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
        lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
        ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
      and trn': trn' = (if trv1 = [] then L else R)
      and lltrv2: ltrv2 =
        lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
      using lltrv1-lltrv2-lnever-L[OF unω Δ r ltr1 isi4 current]
      unfolding ltrv2 by blast
      define ltrv2' where ltrv2': ltrv2' ≡ lltrv2 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
      have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
      unfolding lltrv2 ltrv2' ..
      have nev2: never isSecV trv2 using ω3 unfolding ω3-def by auto
      show ?thesis
      proof(cases trv1 = [])
        case True note trv1 = True
        have [simp]: trn' = trn using L trv1 trn' by auto
        have w1' < w1 using ω3 trv1 unfolding ω3-def by auto
        show ?thesis
        apply(rule TwoFuncPred.sameFM2-selectDelayL)
        apply(rule exI[of - w1']) apply(rule exI[of - nL])
        apply(rule exI[of - trv2]) apply(rule exI[of - []])
        apply(rule exI[of - w2'])
        apply(rule exI[of - ltrv2']) apply(rule exI[of - ltr2])
        apply(rule exI[of - nR])
        apply(intro conjI)
        subgoal by simp subgoal .. subgoal ..
        subgoal by fact subgoal by simp
        subgoal unfolding nL by fact
        subgoal using nev2 by (simp add: never-Nil-filter)
        subgoal apply(rule disjI1)
          apply(rule exI[of - w1']) apply(rule exI[of - w2'])
          apply(rule exI[of - s1']) apply(rule exI[of - s2])
          apply(rule exI[of - ltr1']) apply(rule exI[of - statA])
          apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])

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apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal .. subgoal ..
    subgoal unfolding ltrv2' by simp
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by simp
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Van.reach-validFromS-reach r(4) snoc-eq-iff-butlast)
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$  by auto
  subgoal by fact subgoal by fact subgoal by fact . .
next
case False note trv1 = False
have [simp]: trn' = R using L trv1 trn' by auto
show ?thesis
apply(rule TwoFuncPred.sameFM2-selectLR)
apply(rule exI[of - trv2]) apply(rule exI[of - []])
apply(rule exI[of - w1 ^]) apply(rule exI[of - w2 ^])
apply(rule exI[of - ltrv2 ^]) apply(rule exI[of - ltr2])
apply(rule exI[of - w1]) apply(rule exI[of - w2])
apply(intro conjI)
  subgoal by fact
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal unfolding ltrv2 ..
  subgoal unfolding  $\omega\omega(1)$  by simp
subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by (simp add: never-Nil-filter)
  subgoal apply(rule disjI1)
  apply(rule exI[of - w1 ^]) apply(rule exI[of - w2 ^])
  apply(rule exI[of - s1 ^]) apply(rule exI[of - s2])
  apply(rule exI[of - ltr1 ^]) apply(rule exI[of - statA])
  apply(rule exI[of - sv1 ^]) apply(rule exI[of - sv2 ^])
  apply(rule exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal ..
    subgoal unfolding ltrv2' by simp
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by simp
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Van.reach-validFromS-reach r(4) snoc-eq-iff-butlast)
    subgoal using  $\omega\omega$  by auto

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      subgoal using  $\omega\omega$  by auto
      subgoal using  $\omega\omega$  by auto
      subgoal by fact subgoal by fact subgoal by fact . .
    qed
  next
  case False note current = current False
  obtain  $w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO$  where
   $\chi\chi: \chi\chi s1 ltr1 tr1 s1' s1'' ltr1'$  and
   $\chi\chi': \chi\chi' \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2$ 
   $trv2 sv2'' statOO$ 
  and  $lltrv2: ltrv2 =$ 
   $lappend (l\text{list-of } trv2) (lltrv2 (R, w1', w2', s1'', s1'' \$ ltr1', s2, ltr2, statA, sv1'', sv2'', statOO))$ 

  using  $lltrv1\text{-}lltrv2\text{-not-lnever-L[OF unW } \Delta r ltr1 \text{ isi4 current}]$ 
  unfolding  $ltrv2$  by blast
  define  $ltrv2'$  where  $ltrv2': ltrv2' \equiv lltrv2 (R, w1', w2', s1'', s1'' \$$ 
   $ltr1', s2, ltr2, statA, sv1'', sv2'', statOO)$ 
  have  $ltrv2: ltrv2 = lappend (l\text{list-of } trv2) ltrv2'$ 
  unfolding  $lltrv2 ltrv2'$  ..

  show ?thesis
  apply (rule TwoFuncPred.sameFM2-selectLR)
  apply (rule exI[of -  $trv2$ ]) apply (rule exI[of - []])
  apply (rule exI[of -  $w1'$ ]) apply (rule exI[of -  $w2'$ ])
  apply (rule exI[of -  $ltrv2'$ ]) apply (rule exI[of -  $ltr2$ ])
  apply (rule exI[of -  $w1$ ]) apply (rule exI[of -  $w2$ ])
  apply (intro conjI)
  subgoal by fact
  subgoal unfolding  $nL$  .. subgoal unfolding  $nR$  ..
  subgoal unfolding  $ltrv2$  .. subgoal by simp
  subgoal using  $\chi\chi'$  unfolding  $\chi\chi'\text{-def}$  by (simp add: never-Nil-filter)
  subgoal apply (rule disjI1)
  apply (rule exI[of -  $w1'$ ]) apply (rule exI[of -  $w2'$ ])
  apply (rule exI[of -  $s1''$ ]) apply (rule exI[of -  $s2$ ])
  apply (rule exI[of -  $s1'' \$ ltr1'$ ]) apply (rule exI[of -  $statA$ ])
  apply (rule exI[of -  $sv1''$ ]) apply (rule exI[of -  $sv2''$ ])
  apply (rule exI[of -  $statOO$ ])
  apply (intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding  $ltrv2'$  by simp
  subgoal using  $\chi\chi'$  unfolding  $\chi\chi'\text{-def}$  by simp
  subgoal using  $r(1) \chi\chi$  unfolding  $\chi\chi\text{-def}$ 
  by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
  not-Cons-self2)
  subgoal by fact
  subgoal using  $r(3) \chi\chi'$  unfolding  $\chi\chi'\text{-def}$ 
  by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using  $r(4) \chi\chi'$  unfolding  $\chi\chi'\text{-def}$ 
  by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)

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```

      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of ltr1(3) by blast
      subgoal by fact subgoal by fact subgoal by fact . .
qed
next
case R note trn = R[simp] note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
  case True note current = current True
  obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
   $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
  lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
   $\omega\omega$ :  $\omega\omega$   $\Delta$  w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
  and trn': trn' = (if trv2 = [] then R else L)
  and ltrv2: ltrv2 =
  lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-R[OF un $\omega$   $\Delta$  r ltr2(1,2) isi3 ltr2(3) current]
  unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
  have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding ltrv2 ltrv2' ..
  have nis2:  $\neg$  isSecO s2 using True  $\omega\omega$ (1) by force

show ?thesis
proof(cases trv2 = [])
  case True note trv2 = True
  hence w2' < w2 using  $\omega\omega$  unfolding  $\omega\omega$ -def by auto
  have [simp]: trn' = trn by (simp add: trv2 trn')
  show ?thesis
  apply(rule TwoFuncPred.sameFM2-selectDelayR)
  apply(rule exI[of - w2']) apply(rule exI[of - w2])
  apply(rule exI[of - trv2]) apply(rule exI[of - [s2]])
  apply(rule exI[of - w1'])
  apply(rule exI[of - ltrv2']) apply(rule exI[of - ltr2'])
  apply(rule exI[of - w1])
  apply(intro conjI)
  subgoal by fact
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal unfolding ltrv2 trv2 by simp
  subgoal unfolding  $\omega\omega$ (1) by simp
  subgoal by fact subgoal unfolding trv2 using  $\omega\omega$ -def nis2 by
simp
  subgoal apply(rule disjI1)
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])

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apply(rule exI[of - s1]) apply(rule exI[of - s2'])
apply(rule exI[of - ltr1]) apply(rule exI[of - statA])
apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv2' by simp
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(2) snd-conv)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal by fact subgoal by fact subgoal by fact
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$ 
  using llist-all-lappend-llist-of ltr1(3) by blast . .
next
case False note trv2 = False
show ?thesis
apply(rule TwoFuncPred.sameFM2-selectlappend)
apply(rule exI[of - trv2]) apply(rule exI[of - [s2]])
apply(rule exI[of - trn']) apply(rule exI[of - w1'])
apply(rule exI[of - w2'])
apply(rule exI[of - ltrv2']) apply(rule exI[of - ltr2'])
apply(rule exI[of - trn])
apply(rule exI[of - w1])
apply(rule exI[of - w2])
apply(intro conjI)
  subgoal ..
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrv2 .
  subgoal unfolding  $\omega\omega$ (1) by simp
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal using ltr1(3)  $\omega_4$  unfolding  $\omega_4$ -def
  by (simp add: never-Nil-filter nis2)
  subgoal apply(rule disjI1)
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1]) apply(rule exI[of - s2'])
    apply(rule exI[of - ltr1]) apply(rule exI[of - statA])
    apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
    apply(rule exI[of - statOO])
    apply(intro conjI)

```

```

    subgoal .. subgoal ..
    subgoal unfolding ltrv2' ..
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
    subgoal by fact
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(2) snd-conv)
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
    subgoal by fact subgoal by fact subgoal by fact
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$ 
    using llist-all-lappend-llist-of ltr1(3) by blast . .

qed
next
case False note current = current False
obtain  $w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO$  where
 $\chi\chi: \chi\chi s2 ltr2 tr2 s2' s2'' ltr2'$  and
 $\chi4': \chi4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2$ 
 $trv2 sv2'' statOO$ 
and ltrv2: ltrv2 =
  lappend (llist-of trv2) (lltrv2 (L, w1', w2', s1, ltr1, s2'', s2'' $ ltr2',
statA, sv1'', sv2'', statOO))
  using lltrv1-lltrv2-not-lnever-R[OF un $\omega$   $\Delta$  r ltr2(1,2) isi3 ltr2(3)
current]
  unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (L, w1', w2', s1, ltr1, s2'',
s2'' $ ltr2', statA, sv1'', sv2'', statOO)
  have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding ltrv2 ltrv2' ..
  show ?thesis
  apply(rule TwoFuncPred.sameFM2-selectlappend)
  apply(rule exI[of - trv2]) apply(rule exI[of - tr2 ## s2'])
  apply(rule exI[of - L])
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - ltrv2']) apply(rule exI[of - s2'' $ ltr2'])
  apply(rule exI[of - trn])
  apply(rule exI[of - w1]) apply(rule exI[of - w2])
  apply(intro conjI)
  subgoal .. subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrv2 .
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by simp
  subgoal using  $\chi4'$  unfolding  $\chi4'$ -def by simp
  subgoal by simp
  subgoal using  $\chi4'$  unfolding  $\chi4'$ -def

```



```

    by (simp add: Opt.S.map-filter Van.S.map-filter)
    subgoal apply(rule disjI1)
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1])
    apply(rule exI[of - s2']) apply(rule exI[of - ltr1])
    apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule exI[of
- sv2'])
    apply(rule exI[of - statOO])
    apply(intro conjI)
    subgoal .. subgoal ..
    subgoal unfolding ltrv2' ..
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
    subgoal by fact
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$   $\chi\chi$ -def

    append-is-Nil-conv last-snoc not-Cons-self2 r(2))
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
    by (metis Van.reach-validFromS-reach r(4) snoc-eq-iff-butlast)
    subgoal by fact subgoal by fact subgoal by fact
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
    using llist-all-lappend-llist-of ltr2(3) by blast . .

    qed
  qed
  qed
  qed
  qed
}
thus ?thesis unfolding Van.lS[OF cltrv2] Opt.lS[OF ltr2(2)]
apply- apply(rule TwoFuncPred.sameFM2-lmap-lfilter)
using assms by blast
qed

```

lemma *lA-lltrv1-lltrv2*:

assumes *unw*: *unwindCond* Δ

and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$

and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*

and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1* *lnever* *isIntO* *ltr1*

and *ltr2*: *Opt.lvalidFromS* *s2* *ltr2* *Opt.lcompletedFrom* *s2* *ltr2* *lnever* *isIntO* *ltr2*

shows *Van.lA* (*lltrv1* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*)) =

Van.lA (*lltrv2* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*))

proof –

have *cltrv1*: *Van.lcompletedFrom* *sv1* (*lltrv1* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*))

```

using lcompletedFrom-lltrv1 [OF assms] .
have cltrv2: Van.lcompletedFrom sv2 (lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
using lcompletedFrom-lltrv2 [OF assms] .
{fix nL nR ltrv1 ltrv2
  assume  $\exists$  trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO.
    nL = w1  $\wedge$  nR = w2  $\wedge$ 
    ltrv1 = lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
    ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
     $\Delta \infty$  w1 w2 s1 s2 statA sv1 sv2 statO  $\wedge$ 
    reachO s1  $\wedge$  reachO s2  $\wedge$  reachV sv1  $\wedge$  reachV sv2  $\wedge$ 
    Opt.lvalidFromS s1 ltr1  $\wedge$  Opt.lcompletedFrom s1 ltr1  $\wedge$  lnever isIntO ltr1  $\wedge$ 

    Opt.lvalidFromS s2 ltr2  $\wedge$  Opt.lcompletedFrom s2 ltr2  $\wedge$  lnever isIntO ltr2
  hence TwoFuncPred.sameFM isIntV isIntV getActV getActV nL nR ltrv1 ltrv2
  proof(coinduct rule: TwoFuncPred.sameFM.coinduct[of  $\lambda$ nL nR ltrv1 ltrv2.
     $\exists$  trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO.
      nL = w1  $\wedge$  nR = w2  $\wedge$ 
      ltrv1 = lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
      ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
       $\Delta \infty$  w1 w2 s1 s2 statA sv1 sv2 statO  $\wedge$ 
      reachO s1  $\wedge$  reachO s2  $\wedge$  reachV sv1  $\wedge$  reachV sv2  $\wedge$ 
      Opt.lvalidFromS s1 ltr1  $\wedge$  Opt.lcompletedFrom s1 ltr1  $\wedge$  lnever isIntO ltr1  $\wedge$ 

      Opt.lvalidFromS s2 ltr2  $\wedge$  Opt.lcompletedFrom s2 ltr2  $\wedge$  lnever isIntO ltr2])
  case (2 nL nR ltrv1 ltrv2)
  then obtain trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO
  where nL: nL = w1 and nR: nR = w2
  and ltrv1: ltrv1 = lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and ltrv2: ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and  $\Delta$ :  $\Delta \infty$  w1 w2 s1 s2 statA sv1 sv2 statO
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
  by auto
  have isi3:  $\neg$  isIntO s1 using ltr1
  by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
    llist.pred-inject(2))
    have isi4:  $\neg$  isIntO s2 using ltr2
  by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
    llist.pred-inject(2))

  show ?case proof(cases ltr1 = []  $\wedge$  ltr2 = [])
    case True note ltr14 = True
    hence ltrv1: ltrv1 = [] unfolding ltrv1 by simp
    show ?thesis using ltr14 unfolding ltrv1 ltrv2 apply-apply(rule Two-
    FuncPred.sameFM-selectLNil) by auto
  next
  case False hence ltr14: ltr1  $\neq$  []  $\vee$  ltr2  $\neq$  [] by auto
  show ?thesis proof(cases llength ltr1  $\leq$  Suc 0  $\vee$  llength ltr2  $\leq$  Suc 0)

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    case True note ltr14 = ltr14 True
  hence ltrv1: ltrv1 = [[sv1]] and ltrv2: ltrv2 = [[sv2]] unfolding ltrv1 ltrv2
by auto
  have llength ltr1 = Suc 0 ∨ llength ltr2 = Suc 0
  by (metis Opt.lcompletedFrom-def Suc-ile-eq True
    lfinite-LNil llength-LNil llist-eq-cong ltr1 (2)
    ltr2 (2) nle-le order-le-imp-less-or-eq zero-enat-def zero-order (3))
  hence finalO s1 ∨ finalO s2
  using Opt.lcompletedFrom-singl ltr1 (1) ltr1 (2) ltr2 (1) ltr2 (2) by blast
  hence fs1: finalO s1 ∧ finalO s2
  using Δ r (1) r (2) r (3) r (4) unw unwindCond-def by auto

  have fsv12: finalV sv1 ∧ finalV sv2
  using Δ fs1 r (1) r (2) r (3) r (4) unw unwindCond-final by blast
  have isv12: ¬ isIntV sv1 ∧ ¬ isIntV sv2
  using fsv12 Van.final-not-isInt by blast
show ?thesis unfolding ltrv1 ltrv2 apply (rule TwoFuncPred.sameFM-selectSingl)

  using isv12 by auto
next
case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0
by auto
show ?thesis proof (cases trn)
  case L note current = current L
  show ?thesis
  proof (cases lnever isSecO ltr1)
    case True note current = current True
    obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
      ωω: ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
      lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
      ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
    and trn': trn' = (if trv1 = [] then L else R)
    and lltrv1: ltrv1 =
      lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
    and lltrv2: ltrv2 =
      lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
    using lltrv1-lltrv2-lnever-L[OF unw Δ r ltr1 isi4 current]
    unfolding ltrv1 ltrv2 by blast
    define ltrv1' where ltrv1': ltrv1' ≡ lltrv1 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
    have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
    unfolding lltrv1 ltrv1' ..
    define ltrv2' where ltrv2': ltrv2' ≡ lltrv2 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
    have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
    unfolding lltrv2 ltrv2' ..

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show ?thesis
apply(rule TwoFuncPred.sameFM-selectlappend)
apply(rule exI[of - trv1]) apply(rule exI[of - w1]) apply(rule exI[of -
w1])
apply(rule exI[of - trv2]) apply(rule exI[of - w2]) apply(rule exI[of -
w2])
apply(rule exI[of - ltrv1]) apply(rule exI[of - ltrv2])
apply(intro conjI)
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using lltrv1 .
  subgoal using lltrv2 .
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by simp
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by simp
subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by (simp add: Van.A.map-filter)

subgoal apply(rule disjI1)
apply(rule exI[of - trn]) apply(rule exI[of - w1]) apply(rule exI[of
- w2])
  apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
  apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
  apply(rule exI[of - statA])
  apply(rule exI[of - sv1]) apply(rule exI[of - sv2])
  apply(rule exI[of - statOO])
  apply(intro conjI)
    subgoal .. subgoal ..
    subgoal unfolding ltrv1' ..
    subgoal unfolding ltrv2' ..
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by simp
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$ 
    using llist-all-lappend-llist-of ltr1(3) by blast
    subgoal using  $\omega\omega$  using ltr2(1) by fastforce
    subgoal by fact
    subgoal by fact . .
next
case False note current = current False
obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
XX: XX s1 ltr1 tr1 s1' s1'' ltr1' and

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       $\chi^{3'}$ :  $\chi^{3'} \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2$ 
\Delta r ltr1 isi4 current]
    unfolding ltrv1 ltrv2 by blast
    define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
    have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
    unfolding lltrv1 ltrv1' ..
    define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
    have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
    unfolding lltrv2 ltrv2' ..

    show ?thesis apply(rule TwoFuncPred.sameFM-selectlappend)
    apply(rule exI[of - trv1]) apply(rule exI[of - w1']) apply(rule exI[of -
w1])
    apply(rule exI[of - trv2]) apply(rule exI[of - w2']) apply(rule exI[of -
w2])
    apply(rule exI[of - ltrv1']) apply(rule exI[of - ltrv2'])
    apply(intro conjI)
    subgoal unfolding nL .. subgoal unfolding nR ..
    subgoal using lltrv1 .
    subgoal using lltrv2 .
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by auto
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by auto
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by (simp add: Van.A.map-filter)

    subgoal apply(rule disjI1)
    apply(rule exI[of - R]) apply(rule exI[of - w1']) apply(rule exI[of -
w2'])
    apply(rule exI[of - s1']) apply(rule exI[of - s1'' $ ltr1'])
    apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
    apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule exI[of
- sv2'])

    apply(rule exI[of - statOO])
    apply(intro conjI)
    subgoal ..subgoal ..
    subgoal unfolding ltrv1' ..
    subgoal unfolding ltrv2' ..
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by simp
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$   $\chi\chi$ -def

    append-is-Nil-conv last-snoc not-Cons-self2 r(1))

```

```

      subgoal by fact
      subgoal using  $\chi^3'$  unfolding  $\chi^3'$ -def
        by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
      subgoal using  $\chi^3'$  unfolding  $\chi^3'$ -def
        by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of ltr1(3) by blast
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def using ltr2(1) by fastforce
      subgoal by fact
      subgoal by fact . .
qed
next
case R note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
  case True note current = current True
  obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
 $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
 $\omega_4$ :  $\omega_4 \Delta w1 w2 w1' w2' s1 s2 s2'$  statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
  and trn': trn' = (if trv2 = [] then R else L)
  and ltrv1: ltrv1 =
    lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
  and ltrv2: ltrv2 =
    lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-R[OF un $\omega \Delta r$  ltr2(1,2) isi3 ltr2(3) current]
  unfolding ltrv1 ltrv2 by blast
  define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
  have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..
  define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
  have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding ltrv2 ltrv2' ..
  show ?thesis
  apply(rule TwoFuncPred.sameFM-selectlappend)
  apply(rule exI[of - trv1]) apply(rule exI[of - w1]) apply(rule exI[of -
w1])
  apply(rule exI[of - trv2]) apply(rule exI[of - w2]) apply(rule exI[of -
w2])
  apply(rule exI[of - ltrv1]) apply(rule exI[of - ltrv2])

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apply(intro conjI)
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using lltrv1 .
  subgoal using lltrv2 .
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by (simp add: Van.A.map-filter)
  subgoal apply(rule disjI1)
  apply(rule exI[of - trn1]) apply(rule exI[of - w11]) apply(rule exI[of
- w21])
  apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
  apply(rule exI[of - s21]) apply(rule exI[of - ltr21])
- sv21])
  apply(rule exI[of - statA]) apply(rule exI[of - sv11]) apply(rule exI[of
  apply(rule exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1' ..
  subgoal unfolding ltrv2' ..
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(2) snd-conv)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3))
snoc-eq-iff-butlast)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4))
snoc-eq-iff-butlast)
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto . .
next
case False note current = current False
obtain w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO where
   $\chi\chi$ :  $\chi\chi$  s2 ltr2 tr2 s2' s2'' ltr2' and
   $\chi\chi'$ :  $\chi\chi'$   $\Delta$  w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
  and ltrv1: ltrv1 =
    lappend (llist-of trv1) (lltrv1 (L, w1', w2', s1, ltr1, s2'', s2'' $ ltr2',
statA, sv1'', sv2'', statOO))
  and ltrv2: ltrv2 =
    lappend (llist-of trv2) (lltrv2 (L, w1', w2', s1, ltr1, s2'', s2'' $ ltr2',
statA, sv1'', sv2'', statOO))
  using lltrv1-lltrv2-not-lnever-R[OF un $\omega$   $\Delta$  r ltr2(1,2) isi3 ltr2(3)
current]

```

```

unfolding ltrv1 ltrv2 by blast
define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (L, w1', w2', s1, ltr1, s2'',
s2'' $ ltr2', statA, sv1'', sv2'', statOO)
have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
unfolding ltrv1 ltrv1' ..
define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (L, w1', w2', s1, ltr1, s2'',
s2'' $ ltr2', statA, sv1'', sv2'', statOO)
have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
unfolding ltrv2 ltrv2' ..

show ?thesis
apply(rule TwoFuncPred.sameFM-selectlappend)
apply(rule exI[of - trv1]) apply(rule exI[of - w1']) apply(rule exI[of -
w1])
apply(rule exI[of - trv2]) apply(rule exI[of - w2']) apply(rule exI[of -
w2])

apply(rule exI[of - ltrv1']) apply(rule exI[of - ltrv2'])
apply(intro conjI)
subgoal unfolding nL .. subgoal unfolding nR ..
subgoal using lltrv1 .
subgoal using lltrv2 .
subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by (simp add: Van.A.map-filter)
subgoal apply(rule disjI1)
apply(rule exI[of - L]) apply(rule exI[of - w1']) apply(rule exI[of -
w2'])

apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
apply(rule exI[of - s2''']) apply(rule exI[of - s2'' $ ltr2'])
apply(rule exI[of - statA])
apply(rule exI[of - sv1''']) apply(rule exI[of - sv2''']) apply(rule exI[of -
- statOO])

apply(intro conjI)
subgoal .. subgoal ..
subgoal unfolding ltrv1' ..
subgoal unfolding ltrv2' ..
subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
subgoal by fact
subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$   $\chi\chi$ -def

append-is-Nil-conv last-snoc not-Cons-self2 r(2))
subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
by (metis Simple-Transition-System.reach-validFromS-reach r(3))
snoc-eq-iff-butlast)
subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
by (metis Simple-Transition-System.reach-validFromS-reach r(4))
snoc-eq-iff-butlast)
subgoal by fact

```



```

      subgoal by fact
      subgoal by fact
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of ltr2(3) by blast . .
    qed
  qed
  qed
  qed
  qed
}
thus ?thesis unfolding Van.LA[OF cltrv1] Van.LA[OF cltrv2]
apply- apply(rule TwoFuncPred.sameFM-lmap-lfilter)
using assms by blast
qed

```

```

fun isN :: ('stateO,'stateV) tuple34  $\Rightarrow$  bool
where
isN (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\longleftrightarrow$  ltr1 = []  $\vee$  ltr2 = []

```

```

fun H-T ::
('stateO,'stateV)tuple34  $\Rightarrow$ 
('stateO,'stateV)tuple12  $\times$ 
('stateO,'stateV)tuple34
where
H-T (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) =
(let (tr1,s1',s1'',ltr1',tr2,s2',s2'',ltr2') =
(SOME k. case k of (tr1,s1',s1'',ltr1',tr2,s2',s2'',ltr2')  $\Rightarrow$ 
 $\varphi\varphi$  s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2')
in let (w1',w2',trv1,sv1'',trv2,sv2'',statAA,statOO) =
(SOME k'. case k' of (w1',w2',trv1,sv1'',trv2,sv2'',statAA,statOO)  $\Rightarrow$ 
 $\varphi'$   $\Delta$  w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA statO
sv1 trv1 sv1'' sv2 trv2 sv2'' statOO)
in ((trv1,sv1'',trv2,sv2'',statAA,statOO),
(w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))
)

```

```

declare H-T.simps[simp del]

```

```

definition H  $\equiv$  fst o H-T

```

```

definition T  $\equiv$  snd o H-T

```

```

fun Econd where Econd (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) = (lnever

```

isIntO ltr1)

fun *E* **where** *E* (*w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*) = *f* (*L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*)

definition *F* :: ('stateO,'stateV)tuple34 ⇒ ('stateO,'stateV)tuple12 llist

where *F* ≡ *ccorec-llist isN H Econd E T*

lemma *F-LNil*:

ltr1 = [] ∨ *ltr2* = [] ⇒ *F* (*w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*) = []

unfolding *F-def* **apply**(*subst llist-ccorec(1)*) **by** *auto*

lemma *F-lnever*:

assumes *ltr1* ≠ [] *ltr2* ≠ [] *lnever isIntO ltr1*

shows *F* (*w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*) = *f* (*L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*)

using *assms* **unfolding** *F-def* **apply**(*subst llist-ccorec(2)*)

subgoal unfolding *E.simps lnull-def* **apply**(*rule f-not-LNil*) **by** *auto*

subgoal using *assms* **by** *auto*

subgoal unfolding *Econd.simps* **by** *auto* .

lemma *F-not-lnever*:

assumes *ltr1* ≠ [] *ltr2* ≠ [] ¬ *lnever isIntO ltr1*

shows *F* (*w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*) =

LCons (*H* (*w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*)) (*F* (*T* (*w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*)))

using *assms* **unfolding** *F-def* **apply**(*subst llist-ccorec(2)*)

subgoal unfolding *E.simps lnull-def* **apply**(*rule f-not-LNil*) **by** *auto*

subgoal using *assms* **by** *auto*

subgoal unfolding *Econd.simps* **by** *auto* .

definition *ltrv1* :: ('stateO,'stateV)tuple34 ⇒ 'stateV llist **where**

ltrv1 tp = *lconcat* (*lmap* ($\lambda(trv1,sv1'',trv2,sv2'',statAA,statOO).$ *llist-of trv1*) (*F tp*))

definition *firstHolds1* :: ('stateO,'stateV)tuple34 ⇒ nat **where**

firstHolds1 tp = *firstNC* (*lmap* ($\lambda(trv1,sv1'',trv2,sv2'',statAA,statOO).$ *trv1*) (*F tp*))

definition *ltrv2* :: ('stateO,'stateV)tuple34 ⇒ 'stateV llist **where**

ltrv2 tp = *lconcat* (*lmap* ($\lambda(trv1,sv1'',trv2,sv2'',statAA,statOO).$ *llist-of trv2*) (*F tp*))

definition *firstHolds2* :: ('stateO,'stateV)tuple34 ⇒ nat **where**

firstHolds2 tp = *firstNC* (*lmap* ($\lambda(trv1,sv1'',trv2,sv2'',statAA,statOO).$ *trv2*) (*F tp*))

lemma *ltrv1-ne-imp*:
assumes *ltrv1 tp* $\neq []$
shows $\exists trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO.$ $(trv1, sv1'', trv2, sv2'', statAA, statOO)$
 $\in lset (F\ tp) \wedge$
 $trv1 \neq []$
using *assms unfolding ltrv1-def unfolding lconcat-eq-LNil-iff* **by force**

lemma *ltrv2-ne-imp*:
assumes *ltrv2 tp* $\neq []$
shows $\exists trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO.$ $(trv1, sv1'', trv2, sv2'', statAA, statOO)$
 $\in lset (F\ tp) \wedge$
 $trv2 \neq []$
using *assms unfolding ltrv2-def unfolding lconcat-eq-LNil-iff* **by force**

lemma *ltrv1-LNil[simp]*:
 $ltr1 = [] \vee ltr2 = [] \implies ltrv1 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = []$
unfolding *ltrv1-def F-LNil* **by simp**
lemma *ltrv2-LNil[simp]*:
 $ltr1 = [] \vee ltr2 = [] \implies ltrv2 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = []$
unfolding *ltrv2-def F-LNil* **by simp**

lemma *ltrv1-lnever*:
assumes *ltr1* $\neq []$ *ltr2* $\neq []$ *lnever isIntO ltr1*
shows $ltrv1 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = lltrv1 (L, w1, w2, s1,$
 $ltr1, s2, ltr2, statA, sv1, sv2, statO)$
unfolding *ltrv1-def F-lnever[OF assms]* *lltrv1-def* **..**

lemma *ltrv2-lnever*:
assumes *ltr1* $\neq []$ *ltr2* $\neq []$ *lnever isIntO ltr1*
shows $ltrv2 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = lltrv2 (L, w1, w2, s1,$
 $ltr1, s2, ltr2, statA, sv1, sv2, statO)$
unfolding *ltrv2-def F-lnever[OF assms]* *lltrv2-def* **..**

lemma *H-T-not-lnever*:
assumes *unw: unwindCond Δ*
and $\Delta: \Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
and *r: reachO s1 reachO s2 reachV sv1 reachV sv2*
and *stat: statA = Diff \longrightarrow statO = Diff*
and *ltr1: Opt.validFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1*

and $ltr2$: $Opt.lvalidFromS\ s2\ ltr2\ Opt.lcompletedFrom\ s2\ ltr2$
and l : $\neg\ lnever\ isIntO\ ltr1\ Opt.lA\ ltr1 = Opt.lA\ ltr2$
shows $\exists\ w1'\ w2'\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2'\ trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO$.

$\varphi\varphi\ s1\ ltr1\ s2\ ltr2\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2' \wedge$
 $\varphi'\ \Delta\ w1\ w2\ w1'\ w2'\ statA\ s1\ tr1\ s1'\ s1''\ s2\ tr2\ s2'\ s2''\ statAA\ statO\ sv1\ trv1\ sv1''\ sv2\ trv2\ sv2''\ statOO \wedge$

$H-T\ (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $((trv1, sv1'', trv2, sv2'', statAA, statOO),$
 $(w1', w2', s1'', s1''\ \$\ ltr1', s2'', s2''\ \$\ ltr2', statAA, sv1'', sv2'', statOO))$

proof –

obtain $tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2'$
where $\varphi\varphi$: $\varphi\varphi\ s1\ ltr1\ s2\ ltr2\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2'$
using $isIntO$ - $\varphi\varphi[OF\ ltr1\ ltr2\ l(2,1)]$
by *auto*

define tp **where**

$tp = (SOME\ k.\ case\ k\ of\ (tr1, s1', s1'', ltr1', tr2, s2', s2'', ltr2') \Rightarrow$
 $\varphi\varphi\ s1\ ltr1\ s2\ ltr2\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2')$

have 0 : $case\ tp\ of\ (tr1, s1', s1'', ltr1', tr2, s2', s2'', ltr2') \Rightarrow$
 $\varphi\varphi\ s1\ ltr1\ s2\ ltr2\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2'$

using $\varphi\varphi$ **unfolding** tp -**def** **apply**– **apply**($rule\ someI$ - ex)
apply($rule\ exI$ [of - $(tr1, s1', s1'', ltr1', tr2, s2', s2'', ltr2')$]) **by** *auto*

obtain $tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2'$ **where**

tp : $tp = (tr1, s1', s1'', ltr1', tr2, s2', s2'', ltr2')$ **by**($cases\ tp$, *auto*)

have $\varphi\varphi$: $\varphi\varphi\ s1\ ltr1\ s2\ ltr2\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2'$
using 0 **unfolding** tp **by** *auto*

obtain $w1'\ w2'\ trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO$

where φ' : $\varphi'\ \Delta\ w1\ w2\ w1'\ w2'\ statA\ s1\ tr1\ s1'\ s1''\ s2\ tr2\ s2'\ s2''\ statAA\ statO\ sv1\ trv1\ sv1''\ sv2\ trv2\ sv2''\ statOO$

using $unwindCond$ - ex - $\varphi'[OF\ unw\ \Delta\ r\ stat, of\ tr1\ s1'\ s1''\ tr2\ s2'\ s2'']$
using $\varphi\varphi$ **unfolding** $\varphi\varphi$ -**def** **by** *auto*

define tp' **where**

$tp' = (SOME\ k'.\ case\ k'\ of\ (w1', w2', trv1, sv1'', trv2, sv2'', statAA, statOO) \Rightarrow$
 $\varphi'\ \Delta\ w1\ w2\ w1'\ w2'\ statA\ s1\ tr1\ s1'\ s1''\ s2\ tr2\ s2'\ s2''\ statAA\ statO\ sv1\ trv1\ sv1''\ sv2\ trv2\ sv2''\ statOO)$

have 1 : $case\ tp'\ of\ (w1', w2', trv1, sv1'', trv2, sv2'', statAA, statOO) \Rightarrow$
 $\varphi'\ \Delta\ w1\ w2\ w1'\ w2'\ statA\ s1\ tr1\ s1'\ s1''\ s2\ tr2\ s2'\ s2''\ statAA\ statO\ sv1\ trv1\ sv1''\ sv2\ trv2\ sv2''\ statOO$

using φ' **unfolding** tp' -**def** **apply**– **apply**($rule\ someI$ - ex)
apply($rule\ exI$ [of - $(w1', w2', trv1, sv1'', trv2, sv2'', statAA, statOO)$]) **by** *auto*

```

obtain  $w1' w2' trv1 sv1'' trv2 sv2'' statAA statOO$  where
 $tp'$ :  $tp' = (w1', w2', trv1, sv1'', trv2, sv2'', statAA, statOO)$  by(cases  $tp'$ , auto)

have  $\varphi'$ :  $\varphi' \Delta w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA statO$ 
 $sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$ 
using 1 unfolding  $tp'$  by auto

show ?thesis
apply(rule  $exI[of - w1']$ ) apply(rule  $exI[of - w2']$ )
apply(rule  $exI[of - tr1]$ ) apply(rule  $exI[of - s1']$ ) apply(rule  $exI[of - s1'']$ )
apply(rule  $exI[of - ltr1']$ )
apply(rule  $exI[of - tr2]$ ) apply(rule  $exI[of - s2']$ ) apply(rule  $exI[of - s2'']$ )
apply(rule  $exI[of - ltr2']$ )
apply(rule  $exI[of - trv1]$ ) apply(rule  $exI[of - sv1'']$ ) apply(rule  $exI[of - trv2]$ )
apply(rule  $exI[of - sv2'']$ )
apply(rule  $exI[of - statAA]$ ) apply(rule  $exI[of - statOO]$ )
apply(intro  $conjI$ )
subgoal using  $\varphi\varphi$  .
subgoal using  $\varphi'$  .
subgoal unfolding  $H-T.simps$ 
unfolding  $tp-def[symmetric]$   $tp$  apply  $simp$ 
unfolding  $tp'-def[symmetric]$   $tp'$  by  $simp$  .
qed

lemma  $ltrv1-ltrv2-not-lnever$ :
assumes  $unw$ :  $unwindCond \Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and  $r$ :  $reachO s1 reachO s2 reachV sv1 reachV sv2$ 
and  $stat$ :  $statA = Diff \longrightarrow statO = Diff$ 
and  $ltr1$ :  $Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1$ 
and  $ltr2$ :  $Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2$ 
and  $l$ :  $\neg lnever isIntO ltr1 Opt.lA ltr1 = Opt.lA ltr2$ 
shows  $\exists w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA$ 
 $statOO$ .
 $\varphi\varphi s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' \wedge$ 
 $\varphi' \Delta w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA statO sv1 trv1$ 
 $sv1'' sv2 trv2 sv2'' statOO \wedge$ 
 $ltrv1 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$ 
 $lappend (l\list-of\ trv1) (ltrv1 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2'' \$ ltr2', statAA, sv1'', sv2'', statOO))$ 
 $\wedge$ 
 $ltrv2 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$ 
 $lappend (l\list-of\ trv2) (ltrv2 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2'' \$ ltr2', statAA, sv1'', sv2'', statOO))$ 
proof –
have  $ltr1NE$ :  $ltr1 \neq []$  using  $l(1)$  by auto
hence  $ltr2NE$ :  $ltr2 \neq []$  using  $l(2)$ 
using  $Opt.lcompletedFrom-def\ ltr2(2)$  by blast
show ?thesis
using  $H-T-not-lnever[OF\ assms]$  apply( $elim\ exE$ )

```

subgoal for $w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2''$
 $statAA statOO$
apply(rule $exI[of - w1']$) **apply**(rule $exI[of - w2']$)
apply(rule $exI[of - tr1]$) **apply**(rule $exI[of - s1']$) **apply**(rule $exI[of - s1'']$)
apply(rule $exI[of - ltr1']$)
apply(rule $exI[of - tr2]$) **apply**(rule $exI[of - s2']$) **apply**(rule $exI[of - s2'']$)
apply(rule $exI[of - ltr2']$)
apply(rule $exI[of - trv1]$) **apply**(rule $exI[of - sv1'']$) **apply**(rule $exI[of - trv2]$)
apply(rule $exI[of - sv2'']$)
apply(rule $exI[of - statAA]$) **apply**(rule $exI[of - statOO]$)
apply(intro $conjI$)
subgoal by simp
subgoal by simp
subgoal unfolding $ltrv1-def$ **apply**(subst $F-not-lnever[OF ltr1NE ltr2NE l(1)]$)
unfolding $H-def T-def$ **by simp**
subgoal unfolding $ltrv2-def$ **apply**(subst $F-not-lnever[OF ltr1NE ltr2NE l(1)]$)
unfolding $H-def T-def$ **by simp . .**
qed

lemma $lvalidFromS-ltrv1$:
assumes unw : $unwindCond \Delta$
and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$
and r : $reachO s1 reachO s2 reachV sv1 reachV sv2$
and $stat$: $statA = Diff \longrightarrow statO = Diff$
and $ltr1$: $Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1$
and $ltr2$: $Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2$
and $ltr1_4$: $Opt.lA ltr1 = Opt.lA ltr2$
shows $Van.lvalidFromS sv1 (ltrv1 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO))$
proof–
{fix $n1 sv1 ltrr1$
assume $\exists w1 w2 s1 ltr1 s2 ltr2 statA sv2 statO.$
 $ltrr1 = ltrv1 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$
 $n1 = w1 \wedge$
 $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$
 $reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$
 $(statA = Diff \longrightarrow statO = Diff) \wedge$
 $Opt.lvalidFromS s1 ltr1 \wedge Opt.lcompletedFrom s1 ltr1 \wedge$
 $Opt.lvalidFromS s2 ltr2 \wedge Opt.lcompletedFrom s2 ltr2 \wedge$
 $Opt.lA ltr1 = Opt.lA ltr2$
hence $Van.llvalidFromS n1 sv1 ltrr1$
proof(coinduct rule: $Van.llvalidFromS.coinduct[of \lambda n1 sv1 ltrr1.$
 $\exists w1 w2 s1 ltr1 s2 ltr2 statA sv2 statO.$
 $ltrr1 = ltrv1 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$
 $n1 = w1 \wedge$
 $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$
 $reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$
 $(statA = Diff \longrightarrow statO = Diff) \wedge$

```

    Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧
    Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧
    Opt.lA ltr1 = Opt.lA ltr2])
  case (llvalidFromS n1 sv1 ltrr1)
  then obtain w1 w2 s1 ltr1 s2 ltr2 statA sv2 statO
  where ltrr1: ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and n1: n1 = w1
  and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and stat: statA = Diff → statO = Diff
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
  and A34: Opt.lA ltr1 = Opt.lA ltr2
  by auto

  have current: ltr1 ≠ [] ltr2 ≠ []
  using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

  show ?case proof(cases lnever isIntO ltr1)
    case True note current = current True
    hence lnever isIntO ltr2
    by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
    ltr2(2))
    note ln34 = True this
    have ltrr1: ltrr1 = ltrv1 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
    statO)
    unfolding ltrr1 ltrv1-lnever[OF current] by simp
    show ?thesis apply(rule Van.llvalidFromS-selectlvalidFromS)
    unfolding ltrr1 apply simp
    apply(rule lvalidFromS-ltrv1)
    using ln34 Δ llvalidFromS ln34(2) ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2)
    r(4) unw by auto
  next
    case False note ln3 = False
    hence ln4: ¬ lnever isIntO ltr2
    by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
    ltr2(2))

  have ltr1 ≠ [[s1]] using ln3 ltr1
  using Opt.final-not-isInt by auto
  hence llength ltr1 > Suc 0
  by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(1) enat-0-iff(1)
  linorder-not-less llength-LNil llist-eq-cong ltr1(1) ltr1(2) nle-le not-less-zero)
  hence ¬ finalO s1
  by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(1) eSuc-enat enat-0

  linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr1(1))
  hence nf12: ¬ finalV sv1 ∧ ¬ finalV sv2
  using Δ r(1) r(2) r(3) r(4) unw unwindCond-def by force

```

```

obtain  $w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA$ 
statOO
  where  $\varphi\varphi: \varphi\varphi s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'$ 
  and  $\varphi': \varphi' \Delta w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA$ 
statO  $sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$ 
  and  $ltrr1: ltrr1 =$ 
   $lappend (l\text{list-of } trv1) (ltrv1 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2'' \$ ltr2', statAA, sv1'', sv2'', statOO))$ 

  using  $ltrv1\text{-}ltrv2\text{-not}\text{-lnever}[OF unw \Delta r stat ltr1 ltr2 ln3 A34]$ 
  unfolding  $ltrr1$  by  $blast$ 
  define  $ltrr1'$  where  $ltrr1': ltrr1' = ltrv1 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2''$ 
 $\$ ltr2', statAA, sv1'', sv2'', statOO)$ 
  have  $ltrr1: ltrr1 = lappend (l\text{list-of } trv1) ltrr1'$ 
  unfolding  $ltrr1$   $ltrr1'$  ..
  have  $ne: trv1 \neq [] \vee (trv1 = [] \wedge w1' < w1)$ 
  using  $\varphi'$  unfolding  $\varphi'$ -def  $ltrr1$  by  $simp$ 

  show  $?thesis$  using  $ne$  proof( $elim\ disjE\ conjE$ )
  assume  $trv1: trv1 \neq []$ 
  show  $?thesis$ 
  apply( $rule\ Van.llvalidFromS\text{-select}lappend$ )
  apply( $rule\ exI[of\ \text{-}\ sv1]$ ) apply( $rule\ exI[of\ \text{-}\ trv1]$ )
  apply( $rule\ exI[of\ \text{-}\ sv1'']$ ) apply( $rule\ exI[of\ \text{-}\ w1']$ )
  apply( $rule\ exI[of\ \text{-}\ ltrr1']$ ) apply( $rule\ exI[of\ \text{-}\ w1]$ )
  apply( $intro\ conjI$ )
  subgoal unfolding  $n1$  .. subgoal ..
  subgoal unfolding  $ltrr1$  ..
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def
  by ( $metis\ Van.validS\text{-}append1\ Van.validFromS\text{-}def\ append\text{-is}\text{-}Nil\text{-}conv$ 
 $hd\text{-}append2$ )
  subgoal by  $fact$ 
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def
  by ( $metis\ Simple\text{-}Transition\text{-}System.validFromS\text{-}def\ Van.validS\text{-}validTrans$ 
 $append\text{-is}\text{-}Nil\text{-}conv\ list.sel(1)\ not\text{-}Cons\text{-}self2\ trv1$ )
  subgoal apply( $rule\ disjI1$ )
  apply( $rule\ exI[of\ \text{-}\ w1']$ ) apply( $rule\ exI[of\ \text{-}\ w2']$ )
  apply( $rule\ exI[of\ \text{-}\ s1'']$ ) apply( $rule\ exI[of\ \text{-}\ s1'' \$ ltr1']$ )
  apply( $rule\ exI[of\ \text{-}\ s2'']$ ) apply( $rule\ exI[of\ \text{-}\ s2'' \$ ltr2']$ )
  apply( $rule\ exI[of\ \text{-}\ statAA]$ ) apply( $rule\ exI[of\ \text{-}\ sv2'']$ ) apply( $rule\ exI[of$ 
 $\text{-}\ sv1'']$ )
  apply( $intro\ conjI$ )
  subgoal unfolding  $ltrr1'$  ..
  subgoal ..
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by  $auto$ 
  subgoal using  $r(1)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
  by ( $metis\ Opt.reach\text{-}validFromS\text{-}reach\ append\text{-is}\text{-}Nil\text{-}conv\ last\text{-}snoc$ 
 $not\text{-}Cons\text{-}self2$ )
  subgoal using  $r(2)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ -def

```



```

      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
trv1)
    subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
next
assume trv1[simp]: trv1 = [] and MM': w1' < w1
hence sv1''[simp]: sv1'' = sv1 using  $\varphi'$  unfolding  $\varphi'$ -def by simp
show ?thesis
apply(rule Van.llvalidFromS-selectDelay)
apply(rule exI[of - w1']) apply(rule exI[of - w1])
apply(rule exI[of - sv1'']) apply(rule exI[of - ltrr1'])
apply(intro conjI)
  subgoal unfolding n1 .. subgoal by simp
  subgoal unfolding ltrr1 by simp subgoal by fact
  subgoal apply(rule disjI1)
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1'']) apply(rule exI[of - s1'' $ ltr1'])
    apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
    apply(rule exI[of - statAA]) apply(rule exI[of - sv2'']) apply(rule exI[of
- statOO])
  apply(intro conjI)
  subgoal unfolding ltrr1' ..
  subgoal ..
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
  subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal unfolding sv1'' by fact
  subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
qed

```

```

    qed
  qed
}
thus ?thesis apply-apply(rule Van.llvalidFromS-imp-lvalidFromS)
using assms by blast
qed

```

```

lemma lvalidFromS-ltrv2:
assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff  $\longrightarrow$  statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and ltr14: Opt.lA ltr1 = Opt.lA ltr2
shows Van.lvalidFromS sv2 (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
proof -
  {fix n2 sv2 ltrr2
  assume  $\exists w1\ w2\ s1\ ltr1\ s2\ ltr2\ statA\ sv1\ statO.$ 
    ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
    n2 = w2  $\wedge$ 
     $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$ 
    reachO s1  $\wedge$  reachO s2  $\wedge$  reachV sv1  $\wedge$  reachV sv2  $\wedge$ 
    (statA = Diff  $\longrightarrow$  statO = Diff)  $\wedge$ 
    Opt.lvalidFromS s1 ltr1  $\wedge$  Opt.lcompletedFrom s1 ltr1  $\wedge$ 
    Opt.lvalidFromS s2 ltr2  $\wedge$  Opt.lcompletedFrom s2 ltr2  $\wedge$ 
    Opt.lA ltr1 = Opt.lA ltr2
  hence Van.llvalidFromS n2 sv2 ltrr2
  proof (coinduct rule: Van.llvalidFromS.coinduct[of  $\lambda n2\ sv2\ ltrr2.$ 
     $\exists w1\ w2\ s1\ ltr1\ s2\ ltr2\ statA\ sv1\ statO.$ 
    ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
    n2 = w2  $\wedge$ 
     $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$ 
    reachO s1  $\wedge$  reachO s2  $\wedge$  reachV sv1  $\wedge$  reachV sv2  $\wedge$ 
    (statA = Diff  $\longrightarrow$  statO = Diff)  $\wedge$ 
    Opt.lvalidFromS s1 ltr1  $\wedge$  Opt.lcompletedFrom s1 ltr1  $\wedge$ 
    Opt.lvalidFromS s2 ltr2  $\wedge$  Opt.lcompletedFrom s2 ltr2  $\wedge$ 
    Opt.lA ltr1 = Opt.lA ltr2])
  case (llvalidFromS n2 sv2 ltrr2)
  then obtain w1 w2 s1 ltr1 s2 ltr2 statA sv1 statO
  where ltrr2: ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and n2: n2 = w2
  and  $\Delta$ :  $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ 
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and stat: statA = Diff  $\longrightarrow$  statO = Diff
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
  and A34: Opt.lA ltr1 = Opt.lA ltr2
  by auto
  }

```

```

have current: ltr1 ≠ [] ltr2 ≠ []
using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

show ?case proof(cases lnever isIntO ltr1)
  case True note current = current True
  hence lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))
  note ln34 = True this
  have ltrr2: ltrr2 = ltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
  unfolding ltrr2 ltrv2-lnever[OF current] by simp
  show ?thesis apply(rule Van.lvalidFromS-selectlvalidFromS)
  unfolding ltrr2 apply simp
  apply(rule lvalidFromS-ltrv2)
  using ln34  $\Delta$  lvalidFromS ln34(2) ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2)
r(3) unw by auto
  next
  case False note ln3 = False
  hence ln4:  $\neg$  lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

  have ltr2 ≠ [[s2]] using ln4 ltr2
  using Opt.final-not-isInt by auto
  hence llength ltr2 > Suc 0
  by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(2) enat-0-iff(2)
linorder-not-less llength-LNil llist-eq-cong ltr2(1) ltr2(2) nle-le not-iless0)
  hence  $\neg$  finalO s2
  by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(2) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr2(1))
  hence nf12:  $\neg$  finalV sv1  $\wedge$   $\neg$  finalV sv2
  using  $\Delta$  r(1) r(2) r(3) r(4) unw unwindCond-def by force

  obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO
  where  $\varphi\varphi$ :  $\varphi\varphi$  s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'
  and  $\varphi'$ :  $\varphi'$   $\Delta$  w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
  and ltrr2: ltrr2 =
lappend (llist-of trv2) (ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO)))

  using ltrv1-ltrv2-not-lnever[OF unw  $\Delta$  r stat ltr1 ltr2 ln3 A34]
  unfolding ltrr2 by blast
  define ltrr2' where ltrr2': ltrr2' = ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
 $\$$  ltr2',statAA,sv1'',sv2'',statOO)
  have ltrr2: ltrr2 = lappend (llist-of trv2) ltrr2'

```

```

unfolding lrr2 lrr2' ..
have ne: trv2 ≠ [] ∨ (trv2 = [] ∧ w2' < w2)
using  $\varphi'$  unfolding  $\varphi'$ -def lrr2 by simp

show ?thesis using ne proof(elim disjE conjE)
  assume trv2: trv2 ≠ []
  show ?thesis
  apply(rule Van.llvalidFromS-selectlappend)
  apply(rule exI[of - sv2]) apply(rule exI[of - trv2])
  apply(rule exI[of - sv2'']) apply(rule exI[of - w2'])
  apply(rule exI[of - lrr2']) apply(rule exI[of - w2])
  apply(intro conjI)
  subgoal unfolding n2 .. subgoal ..
  subgoal unfolding lrr2 ..
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.validS-append1 Van.validFromS-def append-is-Nil-conv
hd-append2)
  subgoal by fact
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def
  by (metis Simple-Transition-System.validFromS-def Van.validS-validTrans
append-is-Nil-conv list.distinct(1) list.sel(1) trv2)
  subgoal apply(rule disjI1)
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1'']) apply(rule exI[of - s1'' $ ltr1'])
  apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
  apply(rule exI[of - statAA]) apply(rule exI[of - sv1'']) apply(rule exI[of
- statOO])
  apply(intro conjI)
  subgoal unfolding lrr2' ..
  subgoal ..
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
  subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
next
  assume trv2[simp]: trv2 = [] and MM': w2' < w2

```

```

hence  $sv2''[simp]: sv2'' = sv2$  using  $\varphi'$  unfolding  $\varphi'$ -def by simp
show ?thesis
apply(rule Van.lvalidFromS-selectDelay)
apply(rule exI[of - w2]) apply(rule exI[of - w2])
apply(rule exI[of - sv2'']) apply(rule exI[of - ltrr2'])
apply(intro conjI)
  subgoal unfolding n2 .. subgoal by simp
  subgoal unfolding ltrr2 by simp subgoal by fact
  subgoal apply(rule disjI1)
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1'']) apply(rule exI[of - s1'' $ ltr1'])
  apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
  apply(rule exI[of - statAA]) apply(rule exI[of - sv1'']) apply(rule exI[of
- statOO])
  apply(intro conjI)
  subgoal unfolding ltrr2' ..
  subgoal ..
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
  subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal unfolding sv2'' by fact
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
qed
qed
qed
}
thus ?thesis apply-apply(rule Van.lvalidFromS-imp-lvalidFromS)
using assms by blast
qed

```

```

lemma lcompletedFrom-ltrv1:
assumes unw: unwindCond  $\Delta$ 
and  $\Delta: \Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff  $\longrightarrow$  statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1

```

```

and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and A34: Opt.lA ltr1 = Opt.lA ltr2
shows Van.lcompletedFrom sv1 (ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
proof –
  {fix ltrr1 assume ltrr1: ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
   and lfin: lfinite ltrr1
   hence list-of ltrr1 ≠ [] ∧ finalV (last (list-of ltrr1))
   using assms(2–) proof(induct length (list-of ltrr1) w1
    arbitrary: w2 ltrr1 s1 ltr1 s2 ltr2 statA sv1 sv2 statO
    rule: less2-induct')
   case (less w1 ltrr1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO)
   hence ltrr1: ltrr1 = ltrv1 (w1,w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)
   and lfin: lfinite ltrr1
   and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
   and r: reachO s1 reachO s2 reachV sv1 reachV sv2
   and stat: statA = Diff → statO = Diff
   and ltr1: Opt.lvalidFromS s1 ltr1 lcompletedFromO s1 ltr1
   and ltr2: Opt.lvalidFromS s2 ltr2 lcompletedFromO s2 ltr2
   and A34: Opt.lA ltr1 = Opt.lA ltr2
   by auto

   have current: ltr1 ≠ [] ltr2 ≠ []
   using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

   show ?case proof(cases lnever isIntO ltr1)
     case True note ln3 = True note current = current True
     hence ln4: lnever isIntO ltr2
     by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))
     note ln34 = True this
     have ltrr1: ltrr1 = lltrv1 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
     unfolding ltrr1 ltrv1-lnever[OF current] by simp
     show ?thesis
     using lcompletedFrom-lltrv1[OF unW Δ r ltr1 ln3 ltr2 ln4, of L]
     using lfin[unfolded ltrr1]
     unfolding Van.lcompletedFrom-def ltrr1[symmetric]
     using llist-of-list-of by fastforce
   next
     case False note ln3 = False
     hence ln4: ¬ lnever isIntO ltr2
     by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

   have ltr1 ≠ [s1] using ln3 ltr1
   using Opt.final-not-isInt by auto
   hence llength ltr1 > Suc 0
   by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(1) enat-0-iff(1)
linorder-not-less llength-LNil llist-eq-cong ltr1(1) ltr1(2) nle-le not-less-zero)

```

hence $\neg \text{finalO } s1$
by (*metis Opt.final-def Opt.lvalidFromS-Cons-iff current(1) eSuc-enat enat-0*)

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr1(1)
hence $\text{nf12: } \neg \text{finalV } sv1 \wedge \neg \text{finalV } sv2$
using $\Delta r(1) r(2) r(3) r(4) \text{ unW unwindCond-def by force}$

obtain $w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' \text{statAA}$
statOO

where $\varphi\varphi: \varphi\varphi s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'$
and $\varphi': \varphi' \Delta w1 w2 w1' w2' \text{statA } s1 tr1 s1' s1'' s2 tr2 s2' s2'' \text{statAA}$
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
and $\text{ltrr1: } \text{ltrr1} =$
 $\text{lappend (l\text{list-of } trv1) (ltrv1 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2'' \$ ltr2', \text{statAA}, sv1'', sv2'', \text{statOO}))}$

using *ltrv1-ltrv2-not-lnever[OF unW Δr stat ltr1 ltr2 ln3 A34]*
unfolding *ltrr1 by blast*
define $\text{ltrr1}'$ **where** $\text{ltrr1}' = \text{ltrv1} (w1', w2', s1'', s1'' \$ ltr1', s2'', s2''$
 $\$ \text{ltr2}', \text{statAA}, sv1'', sv2'', \text{statOO})$
have $\text{ltrr1: } \text{ltrr1} = \text{lappend (l\text{list-of } trv1) \text{ltrr1}'$
unfolding *ltrr1 ltrr1' ..*
have $\text{ne: } trv1 \neq [] \vee (trv1 = [] \wedge w1' < w1)$
using φ' **unfolding** φ' -*def* *ltrr1 by simp*

have $\text{lfin': } \text{lfinite } \text{ltrr1}'$
using *lfin ne unfolding ltrr1 by simp*
have $\text{len: } \text{length (l\text{list-of } \text{ltrr1}')} < \text{length (l\text{list-of } \text{ltrr1})} \vee$
 $\text{length (l\text{list-of } \text{ltrr1}')} = \text{length (l\text{list-of } \text{ltrr1})} \wedge w1' < w1$
using $\text{ne } \text{lfin } \text{lfin}'$ **by** (*simp add: list-of-lappend ltrr1*)

have $0: \text{list-of } \text{ltrr1}' \neq [] \wedge \text{finalV (last (l\text{list-of } \text{ltrr1}'))}$
using $\text{len } \text{proof (elim disjE conjE)}$
assume $\text{len: } \text{length (l\text{list-of } \text{ltrr1}')} < \text{length (l\text{list-of } \text{ltrr1})}$
show *?thesis*
apply(*rule less(1)[OF - ltrr1']*)
subgoal by fact subgoal by fact
subgoal using φ' **unfolding** φ' -*def by simp*
subgoal using $r(1) \varphi\varphi$ **unfolding** $\varphi\varphi$ -*def*
by (*metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc*
not-Cons-self2)
subgoal using $r(2) \varphi\varphi$ **unfolding** $\varphi\varphi$ -*def*
by (*metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc*
not-Cons-self2)
subgoal using $r(3) \varphi'$ **unfolding** φ' -*def*
by (*metis Van.reach-validFromS-reach snoc-eq-iff-butlast*)
subgoal using $r(4) \varphi'$ **unfolding** φ' -*def*
by (*metis Van.reach-validFromS-reach snoc-eq-iff-butlast*)
subgoal using φ' **unfolding** φ' -*def by auto*
subgoal using $\varphi\varphi$ **unfolding** $\varphi\varphi$ -*def by simp*

```

      subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp .
next
assume len: length (list-of ltrr1') = length (list-of ltrr1) w1' < w1
show ?thesis
apply (rule less(2)[OF - - ltrr1'])
  subgoal by fact subgoal unfolding len ..
  subgoal by fact
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
  subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp .
qed
show ?thesis unfolding ltrr1 using 0
by (simp add: lfin' list-of-lappend)
qed
qed
}
thus ?thesis unfolding Van.lcompletedFrom-def by auto
qed

```

```

lemma lcompletedFrom-ltrv2:
assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff  $\longrightarrow$  statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and A34: Opt.lA ltr1 = Opt.lA ltr2
shows Van.lcompletedFrom sv2 (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
proof -
  {fix ltrr2 assume ltrr2: ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and lfin: lfinite ltrr2
  hence list-of ltrr2  $\neq$  []  $\wedge$  finalV (last (list-of ltrr2))
  }

```



```

using assms(2-) proof(induct length (list-of ltrr2) w2
  arbitrary: w1 ltrr2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO
  rule: less2-induct')
case (less w2 ltrr2 w1 s1 ltr1 s2 ltr2 statA sv1 sv2 statO)
hence ltrr2: ltrr2 = ltrv2 (w1,w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)
and lfin: lfinite ltrr2
and  $\Delta: \Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff  $\longrightarrow$  statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 lcompletedFromO s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 lcompletedFromO s2 ltr2
and A34: Opt.lA ltr1 = Opt.lA ltr2
by auto

have current: ltr1  $\neq$  [] ltr2  $\neq$  []
using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

show ?case proof(cases lnever isIntO ltr1)
  case True note ln3 = True note current = current True
  hence ln4: lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))
  note ln34 = True this
  have ltrr2: ltrr2 = lltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
  unfolding ltrr2 ltrv2-lnever[OF current] by simp
  show ?thesis
  using lcompletedFrom-lltrv2[OF unW  $\Delta$  r ltr1 ln3 ltr2 ln4, of L]
  using lfin[unfolding ltrr2]
  unfolding Van.lcompletedFrom-def ltrr2[symmetric]
  using llist-of-list-of by fastforce
next
  case False note ln3 = False
  hence ln4:  $\neg$  lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

  have ltr2  $\neq$  [[s2]] using ln4 ltr2
  using Opt.final-not-isInt by auto
  hence llength ltr2 > Suc 0
  by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(2) enat-0-iff(2)
linorder-not-less llength-LNil llist-eq-cong ltr2(1) ltr2(2) nle-le not-less-zero)
  hence  $\neg$  finalO s2
  by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(2) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr2(1))
  hence nf12:  $\neg$  finalV sv1  $\wedge$   $\neg$  finalV sv2
  using  $\Delta$  r(1) r(2) r(3) r(4) unW unwindCond-def by force

```

```

obtain  $w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA$ 
 $statOO$ 
  where  $\varphi\varphi: \varphi\varphi s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'$ 
  and  $\varphi': \varphi' \Delta w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA$ 
 $statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$ 
  and  $ltrr2: ltrr2 =$ 
 $lappend (l\text{list-of } trv2) (ltrv2 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2'' \$ ltr2', statAA, sv1'', sv2'', statOO))$ 

  using  $ltrv1\text{-}ltrv2\text{-not}\text{-lnever}[OF un\Delta r stat ltr1 ltr2 ln3 A34]$ 
  unfolding  $ltrr2$  by  $blast$ 
  define  $ltrr2'$  where  $ltrr2': ltrr2' = ltrv2 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2''$ 
 $\$ ltr2', statAA, sv1'', sv2'', statOO)$ 
  have  $ltrr2: ltrr2 = lappend (l\text{list-of } trv2) ltrr2'$ 
  unfolding  $ltrr2 ltrr2'$  ..
  have  $ne: trv2 \neq [] \vee (trv2 = [] \wedge w2' < w2)$ 
  using  $\varphi'$  unfolding  $\varphi'\text{-def } ltrr2$  by  $simp$ 

  have  $lfin': lfinite ltrr2'$ 
  using  $lfin ne$  unfolding  $ltrr2$  by  $simp$ 
  have  $len: length (l\text{list-of } ltrr2') < length (l\text{list-of } ltrr2) \vee$ 
 $length (l\text{list-of } ltrr2') = length (l\text{list-of } ltrr2) \wedge w2' < w2$ 
  using  $ne lfin lfin'$  by ( $simp add: l\text{list-of-lappend } ltrr2$ )

  have  $0: l\text{list-of } ltrr2' \neq [] \wedge finalV (last (l\text{list-of } ltrr2'))$ 
  using  $len$  proof( $elim disjE conjE$ )
  assume  $len: length (l\text{list-of } ltrr2') < length (l\text{list-of } ltrr2)$ 
  show  $?thesis$ 
  apply( $rule less(1)[OF - ltrr2']$ )
  subgoal by fact subgoal by fact
  subgoal using  $\varphi'$  unfolding  $\varphi'\text{-def}$  by  $simp$ 
  subgoal using  $r(1) \varphi\varphi$  unfolding  $\varphi\varphi\text{-def}$ 
  by ( $metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc$ 
 $not-Cons-self2$ )
  subgoal using  $r(2) \varphi\varphi$  unfolding  $\varphi\varphi\text{-def}$ 
  by ( $metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc$ 
 $not-Cons-self2$ )
  subgoal using  $r(3) \varphi'$  unfolding  $\varphi'\text{-def}$ 
  by ( $metis Van.reach-validFromS-reach snoc-eq-iff-butlast$ )
  subgoal using  $r(4) \varphi'$  unfolding  $\varphi'\text{-def}$ 
  by ( $metis Van.reach-validFromS-reach snoc-eq-iff-butlast$ )
  subgoal using  $\varphi'$  unfolding  $\varphi'\text{-def}$  by  $auto$ 
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi\text{-def}$  by  $simp$ 
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi\text{-def}$  by  $simp$ 
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi\text{-def}$  by  $simp$ 
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi\text{-def}$  by  $simp$ 
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi\text{-def}$  by  $simp$  .
  next
  assume  $len: length (l\text{list-of } ltrr2') = length (l\text{list-of } ltrr2) w2' < w2$ 
  show  $?thesis$ 

```

```

    apply(rule less(2)[OF - - ltrr2])
    subgoal by fact subgoal unfolding len ..
    subgoal by fact
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
    subgoal using  $r(1)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
        by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
        not-Cons-self2)
    subgoal using  $r(2)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
        by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
        not-Cons-self2)
    subgoal using  $r(3)$   $\varphi'$  unfolding  $\varphi'$ -def
        by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using  $r(4)$   $\varphi'$  unfolding  $\varphi'$ -def
        by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp .
    qed
    show ?thesis unfolding ltrr2 using 0
    by (simp add: lfin' list-of-lappend)
  qed
qed
}
thus ?thesis unfolding Van.lcompletedFrom-def by auto
qed

```

lemma *lS-ltrv1-ltr1*:

assumes *unw*: *unwindCond* Δ

and Δ : $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and r : *reachO* $s1$ *reachO* $s2$ *reachV* $sv1$ *reachV* $sv2$

and *stat*: $statA = Diff \longrightarrow statO = Diff$

and *ltr1*: *Opt.lvalidFromS* $s1$ *ltr1* *Opt.lcompletedFrom* $s1$ *ltr1*

and *ltr2*: *Opt.lvalidFromS* $s2$ *ltr2* *Opt.lcompletedFrom* $s2$ *ltr2*

and $A34$: *Opt.lA* *ltr1* = *Opt.lA* *ltr2*

shows *Van.lS* (*ltrv1* ($w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO$)) = *Opt.lS* *ltr1*

proof –

have *cltrv1*: *Van.lcompletedFrom* $sv1$ (*ltrv1* ($w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO$))

using *lcompletedFrom-ltrv1* [*OF* *assms*] .

{**fix** $nL\ nR\ ltrr1\ ltr1$

assume $\exists w1\ w2\ s1\ s2\ ltr2\ statA\ sv1\ sv2\ statO.$

$nL = w1 \wedge nR = w1 \wedge$

$ltrr1 = ltrv1\ (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$

$\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$

$reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$

$(statA = Diff \longrightarrow statO = Diff) \wedge$
 $Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge$
 $Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge$
 $Opt.lA\ ltr1 = Opt.lA\ ltr2$

hence $TwoFuncPred.sameFM\ isSecV\ isSecO\ getSecV\ getSecO\ nL\ nR\ ltrr1\ ltr1$
proof(coinduct rule: $TwoFuncPred.sameFM.coinduct[of\ \lambda nL\ nR\ ltrr1\ ltr1.$
 $\exists\ w1\ w2\ s1\ s2\ ltr2\ statA\ sv1\ sv2\ statO.$
 $nL = w1 \wedge nR = w1 \wedge$
 $ltrr1 = ltrv1\ (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$
 $\Delta \infty\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$
 $reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$
 $(statA = Diff \longrightarrow statO = Diff) \wedge$
 $Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge$
 $Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge$
 $Opt.lA\ ltr1 = Opt.lA\ ltr2]$)

case $(2\ nL\ nR\ ltrr1\ ltr1)$
then obtain $w1\ w2\ s1\ s2\ ltr2\ statA\ sv1\ sv2\ statO$
where $nL: nL = w1$ **and** $nR: nR = w1$
and $ltrr1: ltrr1 = ltrv1\ (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)$
and $\Delta: \Delta \infty\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
and $r: reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$
and $stat: statA = Diff \longrightarrow statO = Diff$
and $ltr1: Opt.lvalidFromS\ s1\ ltr1\ Opt.lcompletedFrom\ s1\ ltr1$
and $ltr2: Opt.lvalidFromS\ s2\ ltr2\ Opt.lcompletedFrom\ s2\ ltr2$
and $A34: Opt.lA\ ltr1 = Opt.lA\ ltr2$
by *auto*

have $current: ltr1 \neq [] \ ltr2 \neq []$
using $ltr1(2)\ ltr2(2)$ **unfolding** $Opt.lcompletedFrom-def$ **by** *auto*

show $?case\ proof(cases\ lnever\ isIntO\ ltr1)$
case $True$ **note** $ln3 = True$ **note** $current = current\ True$
hence $ln4: lnever\ isIntO\ ltr2$
by $(metis\ Opt.lA\ lfiltermap-LNil-never\ lfiltermap-lmap-lfilter\ ltr1(2)\ A34$
 $ltr2(2))$
note $ln34 = True\ this$
have $ltrr1: ltrr1 = lltrv1\ (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,$
 $statO)$
unfolding $ltrr1\ ltrv1-lnever[OF\ current]$ **by** *simp*

have $clltrv1: Van.lcompletedFrom\ sv1\ (lltrv1\ (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO))$
using $lcompletedFrom-lltrv1[OF\ unW\ \Delta\ r\ ltr1\ ln3\ ltr2\ ln4]$.

show $?thesis$ **apply**(rule $TwoFuncPred.sameFM-selectlmap-lfilter$)
unfolding $ltrr1$ **apply** *simp*
using $lS-lltrv1-ltr1[OF\ unW\ \Delta\ r\ ltr1\ ln3\ ltr2\ ln4, of\ L]$
unfolding $Van.lS[OF\ clltrv1]\ Opt.lS[OF\ ltr1(2)]$.

next
case $False$ **note** $ln3 = False$

```

hence ln4:  $\neg$  lnever isIntO ltr2
by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

have ltr1  $\neq$   $[[s1]]$  using ln3 ltr1
using Opt.final-not-isInt by auto
hence llength ltr1  $>$  Suc 0
by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(1) enat-0-iff(1)
linorder-not-less llength-LNil llist-eq-cong ltr1(1) ltr1(2) nle-le not-less-zero)
hence  $\neg$  finalO s1
by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(1) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr1(1))
hence nf12:  $\neg$  finalV sv1  $\wedge$   $\neg$  finalV sv2
using  $\Delta$  r(1) r(2) r(3) r(4) unw unwindCond-def by force

obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO
where  $\varphi\varphi$ :  $\varphi\varphi$  s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'
and  $\varphi'$ :  $\varphi' \Delta$  w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
and ltrr1: ltrr1 =
lappend (llist-of trv1) (ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

using ltrv1-ltrv2-not-lnever[OF unw  $\Delta$  r stat ltr1 ltr2 ln3 A34]
unfolding ltrr1 by blast
define ltrr1' where ltrr1': ltrr1' = ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
$ ltr2',statAA,sv1'',sv2'',statOO)
have ltrr1: ltrr1 = lappend (llist-of trv1) ltrr1'
unfolding ltrr1 ltrr1' ..
have ne1: trv1  $\neq$   $[]$   $\vee$  w1' < w1
using  $\varphi'$  unfolding  $\varphi'$ -def ltrr1 by simp

show ?thesis
apply(rule TwoFuncPred.sameFM-selectlappend)
apply(rule exI[of - trv1]) apply(rule exI[of - w1']) apply(rule exI[of - w1])

apply(rule exI[of - tr1 ## s1']) apply(rule exI[of - w1']) apply(rule exI[of
- w1])
apply(rule exI[of - ltrr1']) apply(rule exI[of - s1'' $ ltr1'])
apply(intro conjI)
subgoal unfolding nL .. subgoal unfolding nR ..
subgoal using ltrr1 .
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
subgoal by fact
subgoal by simp
subgoal using  $\varphi'$  unfolding  $\varphi'$ -def unfolding Van.S.map-filter Opt.S.map-filter
by simp
subgoal apply(rule disjI1)

```

```

apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1''])
apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
apply(rule exI[of - statAA])
apply(rule exI[of - sv1'']) apply(rule exI[of - sv2''])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrr1' ..
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
  subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
  qed
  qed
}
thus ?thesis apply-
unfolding Van.lS[OF cltrv1] Opt.lS[OF ltr1(2)] apply(rule TwoFuncPred.sameFM-lmap-lfilter)
using assms by blast
qed

```

lemma lS-ltrv2-ltr2:

```

assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff  $\longrightarrow$  statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and A34: Opt.lA ltr1 = Opt.lA ltr2
shows Van.lS (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)) = Opt.lS ltr2
proof -
  have cltrv2: Van.lcompletedFrom sv2 (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
using lcompletedFrom-ltrv2[OF assms] .
  {fix nL nR ltrr2 ltr2
  assume  $\exists w1 w2 s1 s2 ltr1 statA sv1 sv2 statO.$ 
    nL = w2  $\wedge$  nR = w2  $\wedge$ 

```

$ltrr2 = ltrv2 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$
 $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$
 $reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$
 $(statA = Diff \longrightarrow statO = Diff) \wedge$
 $Opt.lvalidFromS s1 ltr1 \wedge Opt.lcompletedFrom s1 ltr1 \wedge$
 $Opt.lvalidFromS s2 ltr2 \wedge Opt.lcompletedFrom s2 ltr2 \wedge$
 $Opt.lA ltr1 = Opt.lA ltr2$

hence $TwoFuncPred.sameFM isSecV isSecO getSecV getSecO nL nR ltrr2 ltr2$
proof(coinduct rule: $TwoFuncPred.sameFM.coinduct[of \lambda nL nR ltrr2 ltr2.$
 $\exists w1 w2 s1 s2 ltr1 statA sv1 sv2 statO.$
 $nL = w2 \wedge nR = w2 \wedge$
 $ltrr2 = ltrv2 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$
 $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$
 $reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$
 $(statA = Diff \longrightarrow statO = Diff) \wedge$
 $Opt.lvalidFromS s1 ltr1 \wedge Opt.lcompletedFrom s1 ltr1 \wedge$
 $Opt.lvalidFromS s2 ltr2 \wedge Opt.lcompletedFrom s2 ltr2 \wedge$
 $Opt.lA ltr1 = Opt.lA ltr2]$)

case ($2 nL nR ltrr2 ltr2$)
then obtain $w1 w2 s1 s2 ltr1 statA sv1 sv2 statO$
where $nL: nL = w2$ **and** $nR: nR = w2$
and $ltrr2: ltrr2 = ltrv2 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)$
and $\Delta: \Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$
and $r: reachO s1 reachO s2 reachV sv1 reachV sv2$
and $stat: statA = Diff \longrightarrow statO = Diff$
and $ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1$
and $ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2$
and $A34: Opt.lA ltr1 = Opt.lA ltr2$
by auto

have current: $ltr1 \neq []$ $ltr2 \neq []$
using $ltr1(2)$ $ltr2(2)$ **unfolding** $Opt.lcompletedFrom-def$ **by auto**

show $?case$ **proof**(cases $lnever isIntO ltr1$)
case $True$ **note** $ln3 = True$ **note** $current = current True$
hence $ln4: lnever isIntO ltr2$
by ($metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34$
 $ltr2(2)$)
note $ln34 = True$ **this**
have $ltrr2: ltrr2 = lltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,$
 $statO)$
unfolding $ltrr2 ltrv2-lnever[OF current]$ **by simp**

have $cllrv2: Van.lcompletedFrom sv2 (lltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO))$
using $lcompletedFrom-lltrv2[OF unW \Delta r ltr1 ln3 ltr2 ln4]$.

show $?thesis$ **apply**(rule $TwoFuncPred.sameFM-selectlmap-lfilter$)
unfolding $ltrr2$ **apply** $simp$
using $lS-lltrv2-ltr2[OF unW \Delta r ltr1 ln3 ltr2 ln4, of L]$

```

unfolding Van.lS[OF cltrv2] Opt.lS[OF ltr2(2)] .
next
  case False note ln3 = False
  hence ln4: ¬ lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

  have ltr2 ≠ [[s2]] using ln4 ltr2
  using Opt.final-not-isInt by auto
  hence llength ltr2 > Suc 0
  by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(2) enat-0-iff(2)
linorder-not-less llength-LNil llist-eq-cong ltr2(1) ltr2(2) nle-le not-less-zero)
  hence ¬ finalO s2
  by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(2) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr2(1))
  hence nf12: ¬ finalV sv1 ∧ ¬ finalV sv2
  using  $\Delta$  r(1) r(2) r(3) r(4) unW unwindCond-def by force

  obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO
  where  $\varphi\varphi$ :  $\varphi\varphi$  s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'
  and  $\varphi'$ :  $\varphi' \Delta$  w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
  and ltrr2: ltrr2 =
lappend (llist-of trv2) (ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

  using ltrv1-ltrv2-not-lnever[OF unW  $\Delta$  r stat ltr1 ltr2 ln3 A34]
  unfolding ltrr2 by blast
  define ltrr2' where ltrr2' = ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
$ ltr2',statAA,sv1'',sv2'',statOO)
  have ltrr1: ltrr2 = lappend (llist-of trv2) ltrr2'
  unfolding ltrr2 ltrr2' ..
  have ne2: trv2 ≠ [] ∨ w2' < w2
  using  $\varphi'$  unfolding  $\varphi'$ -def ltrr2 by simp

  show ?thesis
  apply(rule TwoFuncPred.sameFM-selectlappend)
  apply(rule exI[of - trv2]) apply(rule exI[of - w2']) apply(rule exI[of - w2])

  apply(rule exI[of - tr2 ## s2']) apply(rule exI[of - w2']) apply(rule exI[of
- w2])
  apply(rule exI[of - ltrr2']) apply(rule exI[of - s2'' $ ltr2'])
  apply(intro conjI)
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrr1 .
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal by fact
  subgoal by simp

```



```

    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def unfolding Van.S.map-filter Opt.S.map-filter
  by simp
    subgoal apply (rule disjI1)
      apply (rule exI[of - w1] ) apply (rule exI[of - w2] )
      apply (rule exI[of - s1] )
      apply (rule exI[of - s2] ) apply (rule exI[of - s1'' $ ltr1] )
      apply (rule exI[of - statAA] )
      apply (rule exI[of - sv1] ) apply (rule exI[of - sv2] )
      apply (rule exI[of - statOO] )
      apply (intro conjI)
        subgoal .. subgoal ..
          subgoal unfolding ltr2' ..
            subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
              subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
                by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
              subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
                by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
              subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
                by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
              subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
                by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
              subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
              subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
              subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
              subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
              subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
              subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
            qed
          qed
        }
      thus ?thesis apply-
        unfolding Van.lS[OF cltrv2] Opt.lS[OF ltr2(2)] apply (rule TwoFuncPred.sameFM-lmap-lfilter)
        using assms by blast
    qed

```

lemma *lA-ltrv1-ltrv2*:

assumes *unw*: *unwindCond* Δ

and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$

and *r*: *reachO* *s1 reachO s2 reachV sv1 reachV sv2*

and *stat*: *statA* = *Diff* \longrightarrow *statO* = *Diff*

and *ltr1*: *Opt.lvalidFromS* *s1 ltr1 Opt.lcompletedFrom* *s1 ltr1*

and *ltr2*: *Opt.lvalidFromS* *s2 ltr2 Opt.lcompletedFrom* *s2 ltr2*

and *A34*: *Opt.lA* *ltr1* = *Opt.lA* *ltr2*

shows *Van.lA* (*ltrv1* (*w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO*)) =
Van.lA (*ltrv2* (*w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO*))

```

proof–
  have cltrv1: Van.lcompletedFrom sv1 (ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  using lcompletedFrom-ltrv1[OF assms] .
  have cltrv2: Van.lcompletedFrom sv2 (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  using lcompletedFrom-ltrv2[OF assms] .
  {fix nL nR ltrr1 ltrr2
  assume  $\exists w1\ w2\ s1\ ltr1\ s2\ ltr2\ statA\ sv1\ sv2\ statO.$ 
     $nL = w1 \wedge nR = w2 \wedge$ 
     $ltrr1 = ltrv1\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) \wedge$ 
     $ltrr2 = ltrv2\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) \wedge$ 
     $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$ 
     $reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$ 
     $(statA = Diff \longrightarrow statO = Diff) \wedge$ 
     $Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge$ 
     $Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge$ 
     $Opt.lA\ ltr1 = Opt.lA\ ltr2$ 
  hence TwoFuncPred.sameFM isIntV isIntV getActV getActV nL nR ltrr1 ltrr2
  proof(coinduct rule: TwoFuncPred.sameFM.coinduct[of  $\lambda nL\ nR\ ltrr1\ ltrr2.$ 
     $\exists w1\ w2\ s1\ ltr1\ s2\ ltr2\ statA\ sv1\ sv2\ statO.$ 
     $nL = w1 \wedge nR = w2 \wedge$ 
     $ltrr1 = ltrv1\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) \wedge$ 
     $ltrr2 = ltrv2\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) \wedge$ 
     $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$ 
     $reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$ 
     $(statA = Diff \longrightarrow statO = Diff) \wedge$ 
     $Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge$ 
     $Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge$ 
     $Opt.lA\ ltr1 = Opt.lA\ ltr2]$ )
  case ( $2\ nL\ nR\ ltrr1\ ltrr2$ )
  then obtain  $w1\ w2\ s1\ ltr1\ s2\ ltr2\ statA\ sv1\ sv2\ statO$ 
  where  $nL: nL = w1$  and  $nR: nR = w2$ 
  and  $ltrr1: ltrr1 = ltrv1\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)$ 
  and  $ltrr2: ltrr2 = ltrv2\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)$ 
  and  $\Delta: \Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ 
  and  $r: reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$ 
  and  $stat: statA = Diff \longrightarrow statO = Diff$ 
  and  $ltr1: Opt.lvalidFromS\ s1\ ltr1\ Opt.lcompletedFrom\ s1\ ltr1$ 
  and  $ltr2: Opt.lvalidFromS\ s2\ ltr2\ Opt.lcompletedFrom\ s2\ ltr2$ 
  and  $A34: Opt.lA\ ltr1 = Opt.lA\ ltr2$ 
  by auto

  have current:  $ltr1 \neq [] \wedge ltr2 \neq []$ 
  using  $ltr1(2)\ ltr2(2)$  unfolding Opt.lcompletedFrom-def by auto

  show ?case proof(cases lnever isIntO ltr1)
    case True note  $ln3 = True$  note  $current = current\ True$ 
    hence  $ln4: lnever\ isIntO\ ltr2$ 
    by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
     $ltr2(2)$ )

```

```

note  $ln34 = True$  this
  have  $ltrr1$ :  $ltrr1 = lltrv1 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,$ 
statO)
  unfolding  $ltrr1$   $ltrv1$ - $lnever$ [OF current] by simp
  have  $ltrr2$ :  $ltrr2 = lltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,$ 
statO)
  unfolding  $ltrr2$   $ltrv2$ - $lnever$ [OF current] by simp

have  $clltrv1$ :  $Van.lcompletedFrom sv1 (lltrv1 (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))$ 
  using  $lcompletedFrom$ - $lltrv1$ [OF un $\Delta$  r ltr1 ln3 ltr2 ln4].
have  $clltrv2$ :  $Van.lcompletedFrom sv2 (lltrv2 (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))$ 
  using  $lcompletedFrom$ - $lltrv2$ [OF un $\Delta$  r ltr1 ln3 ltr2 ln4].

show ?thesis apply(rule TwoFuncPred.sameFM-selectlmap-lfilter)
unfolding  $ltrr1$   $ltrr2$  apply simp
using  $lA$ - $lltrv1$ - $lltrv2$ [OF un $\Delta$  r ltr1 ln3 ltr2 ln4, of L]
unfolding  $Van.lA$ [OF clltrv1]  $Van.lA$ [OF clltrv2].
next
case False note  $ln3 = False$ 
hence  $ln4$ :  $\neg lnever$  isIntO  $ltr2$ 
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

have  $ltr1 \neq [[s1]]$  using  $ln3$   $ltr1$ 
using  $Opt.final-not-isInt$  by auto
hence  $llength$   $ltr1 > Suc 0$ 
by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(1) enat-0-iff(1)
linorder-not-less llength-LNil llist-eq-cong ltr1(1) ltr1(2) nle-le not-less-zero)
hence  $\neg finalO s1$ 
by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(1) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr1(1))
hence  $nf12$ :  $\neg finalV sv1 \wedge \neg finalV sv2$ 
using  $\Delta r(1) r(2) r(3) r(4) un $\Delta$  unwindCond-def$  by force

obtain  $w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA$ 
statOO
  where  $\varphi\varphi$ :  $\varphi\varphi s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'$ 
  and  $\varphi'$ :  $\varphi' \Delta w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA$ 
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
  and  $ltrr1$ :  $ltrr1 =$ 
lappend (l $list$ -of  $trv1$ ) (ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

and  $ltrr2$ :  $ltrr2 =$ 
lappend (l $list$ -of  $trv2$ ) (ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

using  $ltrv1$ - $ltrv2$ - $not-lnever$ [OF un $\Delta$  r stat ltr1 ltr2 ln3 A34]
unfolding  $ltrr1$   $ltrr2$  by blast
define  $ltrr1'$  where  $ltrr1'$ :  $ltrr1' = ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''$ 

```

```

$ ltr2',statAA,sv1'',sv2'',statOO)
  have ltrr1: ltrr1 = lappend (llist-of trv1) ltrr1'
  unfolding ltrr1 ltrr1' ..
  have ne1: trv1 ≠ [] ∨ w1' < w1
  using φ' unfolding φ'-def ltrr1 by simp
  define ltrr2' where ltrr2': ltrr2' = ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
$ ltr2',statAA,sv1'',sv2'',statOO)
  have ltrr2: ltrr2 = lappend (llist-of trv2) ltrr2'
  unfolding ltrr2 ltrr2' ..
  have ne2: trv2 ≠ [] ∨ w2' < w2
  using φ' unfolding φ'-def ltrr1 by simp

show ?thesis
apply(rule TwoFuncPred.sameFM-selectlappend)
apply(rule exI[of - trv1]) apply(rule exI[of - w1']) apply(rule exI[of - w1])

apply(rule exI[of - trv2]) apply(rule exI[of - w2']) apply(rule exI[of - w2])

apply(rule exI[of - ltrr1']) apply(rule exI[of - ltrr2'])
apply(intro conjI)
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrr1 .
  subgoal using ltrr2 .
  subgoal by fact subgoal by fact
  subgoal using φ' unfolding φ'-def unfolding Van.A.map-filter by simp
  subgoal apply(rule disjI1)
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1'']) apply(rule exI[of - s1'' $ ltr1'])
    apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
    apply(rule exI[of - statAA])
    apply(rule exI[of - sv1'']) apply(rule exI[of - sv2''])
    apply(rule exI[of - statOO])
    apply(intro conjI)
    subgoal .. subgoal ..
    subgoal unfolding ltrr1' ..
    subgoal unfolding ltrr2' ..
    subgoal using φ' unfolding φ'-def by simp
    subgoal using r(1) φφ unfolding φφ-def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(2) φφ unfolding φφ-def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(3) φ' unfolding φ'-def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using r(4) φ' unfolding φ'-def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using φ' unfolding φ'-def by simp
    subgoal using r(2) φφ unfolding φφ-def by simp

```

```

      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
    qed
  qed
}
thus ?thesis apply-
unfolding Van.lA[OF cltrv1] Van.lA[OF cltrv2] apply(rule TwoFuncPred.sameFM-lmap-lfilter)
using assms by blast
qed

```

lemma *lO-ltrv1-ltrv2*:

```

assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff  $\longrightarrow$  statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and A34: Opt.lA ltr1 = Opt.lA ltr2
and O12: Van.lO (ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)) =
          Van.lO (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
and stO: statO = Eq
shows Opt.lO ltr1 = Opt.lO ltr2
proof-
  have cltrv1: Van.lcompletedFrom sv1 (ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
    using lcompletedFrom-ltrv1[OF assms(1-12)] .
  have cltrv2: Van.lcompletedFrom sv2 (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
    using lcompletedFrom-ltrv2[OF assms(1-12)] .
  {fix nL nR ltr1 ltr2
   assume  $\exists ltrr1 ltrr2 w1 w2 s1 s2 statA sv1 sv2 statO$ .
     ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
     ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
      $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$ 
     reachO s1  $\wedge$  reachO s2  $\wedge$  reachV sv1  $\wedge$  reachV sv2  $\wedge$ 
     (statA = Diff  $\longrightarrow$  statO = Diff)  $\wedge$ 
     Opt.lvalidFromS s1 ltr1  $\wedge$  Opt.lcompletedFrom s1 ltr1  $\wedge$ 
     Opt.lvalidFromS s2 ltr2  $\wedge$  Opt.lcompletedFrom s2 ltr2  $\wedge$ 
     Opt.lA ltr1 = Opt.lA ltr2  $\wedge$ 
     Van.lO ltrr1 = Van.lO ltrr2  $\wedge$ 
     statO = Eq
   }
  hence TwoFuncPred.sameFM isIntO isIntO getObsO getObsO nL nR ltr1 ltr2
  proof(coinduct rule: TwoFuncPred.sameFM.coinduct[of  $\lambda nL nR ltr1 ltr2$ .
     $\exists ltrr1 ltrr2 w1 w2 s1 s2 statA sv1 sv2 statO$ .
      ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
      ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
       $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$ 

```

```

    reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
    (statA = Diff → statO = Diff) ∧
    Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧
    Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧
    Opt.lA ltr1 = Opt.lA ltr2 ∧
    Van.lO ltrr1 = Van.lO ltrr2 ∧
    statO = Eq])
  case (2 nL nR ltr1 ltr2)
  then obtain ltrr1 ltrr2 w1 w2 s1 s2 statA sv1 sv2 statO where
    ltrr11: ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and ltrr22: ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and stat: statA = Diff → statO = Diff
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
  and A34: Opt.lA ltr1 = Opt.lA ltr2
  and O12: Van.lO ltrr1 = Van.lO ltrr2
  and stO: statO = Eq
  by auto

  have stA: statA = Eq using stat stO
  using status.exhaust by blast

  have current: ltr1 ≠ [] ltr2 ≠ []
  using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

  show ?case proof (cases lnever isIntO ltr1)
    case True note ln3 = True note current = current True
    hence ln4: lnever isIntO ltr2
    by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
    ltr2(2))
    note ln34 = True this
    have ltrr1: ltrr1 = lltrv1 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
    statO)
    unfolding ltrr11 ltrv1-lnever[OF current] by simp
    have ltrr2: ltrr2 = lltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
    statO)
    unfolding ltrr22 ltrv2-lnever[OF current] by simp

  have cltrv1: Van.lcompletedFrom sv1 (lltrv1 (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  using lcompletedFrom-lltrv1[OF unW Δ r ltr1 ln3 ltr2 ln4] .
  have cltrv2: Van.lcompletedFrom sv2 (lltrv2 (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  using lcompletedFrom-lltrv2[OF unW Δ r ltr1 ln3 ltr2 ln4] .

  show ?thesis apply (rule TwoFuncPred.sameFM-selectlmap-lfilter)
  unfolding Opt.lO[OF ltr1(2)] Opt.lO[OF ltr2(2)]
  by (metis ln3 ln4 lnever-LNil-lfilter')

```

```

next
  case False note ln3 = False
  hence ln4:  $\neg$  lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

  have ltr1  $\neq$   $[[s1]]$  using ln3 ltr1
  using Opt.final-not-isInt by auto
  hence llength ltr1 > Suc 0
  by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(1) enat-0-iff(1)
linorder-not-less llength-LNil llist-eq-cong ltr1(1) ltr1(2) nle-le not-less-zero)
  hence  $\neg$  finalO s1
  by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(1) eSuc-enat enat-0

  linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr1(1))
  hence nf12:  $\neg$  finalV sv1  $\wedge$   $\neg$  finalV sv2
  using  $\Delta$  r(1) r(2) r(3) r(4) unw unwindCond-def by force

  obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO
  where  $\varphi\varphi$ :  $\varphi\varphi$  s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'
  and  $\varphi'$ :  $\varphi'$   $\Delta$  w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
  and ltrr1: ltrr1 =
  lappend (llist-of trv1) (ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

  and ltrr2: ltrr2 =
  lappend (llist-of trv2) (ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

  using ltrv1-ltrv2-not-lnever[OF unw  $\Delta$  r stat ltr1 ltr2 ln3 A34]
  unfolding ltrr11 ltrr22 by blast
  define ltrr1' where ltrr1': ltrr1' = ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
   $\$$  ltr2',statAA,sv1'',sv2'',statOO)
  have ltrr1: ltrr1 = lappend (llist-of trv1) ltrr1'
  unfolding ltrr1 ltrr1' ..
  have ne1: trv1  $\neq$   $[]$   $\vee$  w1' < w1
  using  $\varphi'$  unfolding  $\varphi'$ -def ltrr1 by simp
  define ltrr2' where ltrr2': ltrr2' = ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
   $\$$  ltr2',statAA,sv1'',sv2'',statOO)
  have ltrr2: ltrr2 = lappend (llist-of trv2) ltrr2'
  unfolding ltrr2 ltrr2' ..
  have ne2: trv2  $\neq$   $[]$   $\vee$  w2' < w2
  using  $\varphi'$  unfolding  $\varphi'$ -def ltrr1 by simp

  have ltr1-eq: ltr1 = lappend (llist-of (tr1 ## s1')) (s1'' $ ltr1')
  and ltr2-eq: ltr2 = lappend (llist-of (tr2 ## s2')) (s2'' $ ltr2') using  $\varphi\varphi$ 
unfolding  $\varphi\varphi$ -def by auto

  have sst: statOO = Diff  $\longleftrightarrow$  Van.O (trv1 ## sv1'')  $\neq$  Van.O (trv2 ##

```

```

sv2'')
  statA = Eq  $\implies$  statAA = Diff  $\longleftrightarrow$  Opt.O ((tr1 ## s1') ## s1'')  $\neq$  Opt.O
((tr2 ## s2') ## s2'')
  statO = Diff  $\implies$  statOO = Diff
  statAA = Diff  $\implies$  statOO = Diff
  using  $\varphi'$  stO unfolding  $\varphi'$ -def by auto

have Atrv12': Van.A (trv1 ## sv1'') = Van.A (trv2 ## sv2'')
using  $\varphi'$  unfolding  $\varphi'$ -def by auto

have  $\Delta'$ :  $\Delta \infty w1' w2' s1'' s2''$  statAA sv1'' sv2'' statOO
using  $\varphi'$  unfolding  $\varphi'$ -def by auto

have vltrv1: Van.lvalidFromS sv1 ltrr1
unfolding ltrr11 using lvalidFromS-ltrv1
using A34  $\Delta$  ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2) r(3) r(4) stat unw
by blast
have cltrv1: Van.lcompletedFrom sv1 ltrr1
unfolding ltrr11 using lcompletedFrom-ltrv1
using A34  $\Delta$  ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2) r(3) r(4) stat unw
by blast

have vltrv2: Van.lvalidFromS sv2 ltrr2
unfolding ltrr22 using lvalidFromS-ltrv2
using A34  $\Delta$  ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2) r(3) r(4) stat unw
by blast
have cltrv2: Van.lcompletedFrom sv2 ltrr2
unfolding ltrr22 using lcompletedFrom-ltrv2
using A34  $\Delta$  ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2) r(3) r(4) stat unw
by blast

have Oltrr1: Van.lO ltrr1 = lmap getObsV (lfilter isIntV ltrr1)
using Van.lO[OF cltrv1] .
have Oltrr2: Van.lO ltrr2 = lmap getObsV (lfilter isIntV ltrr2)
using Van.lO[OF cltrv2] .

have cltrv1': Van.lcompletedFrom (lastt sv1 trv1) ltrr1'
using cltrv1 unfolding ltrr1 Van.lcompletedFrom-def Van.final-def apply
simp
using  $\varphi'$  unfolding  $\varphi'$ -def
by (metis Van.validS-validTrans Van.validFromS-def lappend-LNil2 last-appendR
lfinite-llist-of list-of-lappend
list-of-llist-of llist-of.simps(1) snoc-eq-iff-butlast)

have cltrv2': Van.lcompletedFrom (lastt sv2 trv2) ltrr2'
using cltrv2 unfolding ltrr2 Van.lcompletedFrom-def Van.final-def apply
simp
using  $\varphi'$  unfolding  $\varphi'$ -def
by (metis Van.validS-validTrans Van.validFromS-def lappend-LNil2 last-appendR

```



```

lfinite-llist-of list-of-lappend
  list-of-llist-of llist-of.simps(1) snoc-eq-iff-butlast

  have Oltrr1': Van.lO ltrr1' = lmap getObsV (lfilter isIntV ltrr1')
  using Van.lO[OF cltrv1'] .
  have Oltrr2': Van.lO ltrr2' = lmap getObsV (lfilter isIntV ltrr2')
  using Van.lO[OF cltrv2'] .

  have Van.O (trv1 ## sv1'') = Van.O (trv2 ## sv2'') ∧
    Van.lO ltrr1' = Van.lO ltrr2'
  using Atrv12' O12 Oltrr1' Oltrr2' unfolding Oltrr1 Oltrr2 unfolding ltrr1
ltrr2
  unfolding Van.A.map-filter Van.O.map-filter Van.lO.lmap-lfilter apply simp

  unfolding lmap-lappend-distrib apply simp apply(subst (asm) lappend-llist-of-inj)

  using map-eq-imp-length-eq by auto
  hence O12'': Van.O (trv1 ## sv1'') = Van.O (trv2 ## sv2'')
  and O12': Van.lO ltrr1' = Van.lO ltrr2' by auto

  have stOO: statOO = Eq using O12'' sst by(cases statOO, auto)

  have O34': Opt.O ((tr1 ## s1') ## s1'') = Opt.O ((tr2 ## s2') ##
s2'')
  using stOO sst(2) sst(4) stA by blast

  hence s14': getObsO s1' = getObsO s2'
  using φφ unfolding φφ-def Opt.O.map-filter by (simp add: never-Nil-filter)

  show ?thesis
  apply(rule TwoFuncPred.sameFM-selectlappend)
  apply(rule exI[of - tr1 ## s1']) apply(rule exI[of - undefined]) apply(rule
exI[of - nL])
  apply(rule exI[of - tr2 ## s2']) apply(rule exI[of - undefined]) apply(rule
exI[of - nR])
  apply(rule exI[of - s1'' $ ltr1']) apply(rule exI[of - s2'' $ ltr2'])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal by fact subgoal by fact
  subgoal by simp subgoal by simp
  subgoal using φφ unfolding φφ-def unfolding Opt.O.map-filter
  by (simp add: s14' never-Nil-filter)
  subgoal apply(rule disjI1)
    apply(rule exI[of - ltrr1']) apply(rule exI[of - ltrr2'])
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1'']) apply(rule exI[of - s2''])
    apply(rule exI[of - statAA])
    apply(rule exI[of - sv1'']) apply(rule exI[of - sv2''])
    apply(rule exI[of - statOO])

```

```

    apply(intro conjI)
      subgoal unfolding ltrr1' ..
      subgoal unfolding ltrr2' ..
      subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
      subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
        by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
        by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
      subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
        by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
      subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
        by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
      subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal by fact
      subgoal by fact . .

    qed
  qed
}
thus ?thesis
  unfolding Opt.lO[OF ltr1(2)] Opt.lO[OF ltr2(2)] apply(rule TwoFuncPred.sameFM-lmap-lfilter[where
wL = undefined and wR = undefined])
  using assms by blast
qed
end

```

```

theorem unwind-lrsecure:
assumes init: initCond  $\Delta$  and unwind: unwindCond  $\Delta$ 
shows lrsecure
unfolding lrsecure-def2 proof clarify
  fix s1 tr1 s2 tr2
  assume 3: istateO s1 Opt.lvalidFromS s1 tr1 lcompletedFromO s1 tr1
  and 4: istateO s2 Opt.lvalidFromS s2 tr2 lcompletedFromO s2 tr2
  and A34: Opt.lA tr1 = Opt.lA tr2 and O34: Opt.lO tr1  $\neq$  Opt.lO tr2
  obtain sv1 sv2 where
  isv12: istateV sv1 istateV sv2 and c12: corrState sv1 s1 corrState sv2 s2
  and  $\Delta$ :  $\Delta \infty \infty \infty$  s1 s2 Eq sv1 sv2 Eq
  using init 3 4 unfolding initCond-def by blast
  have r: reachV sv1 reachV sv2 reachO s1 reachO s2

```

by (*auto simp: Van.Istate isv12 Opt.Istate 3 4*)
note $all = 3\ 4\ A34\ isv12\ \Delta\ unwind\ r$
show $\exists sv1\ trv1\ sv2\ trv2.$
 $istateV\ sv1 \wedge istateV\ sv2 \wedge corrState\ sv1\ s1 \wedge corrState\ sv2\ s2 \wedge$
 $Van.lvalidFromS\ sv1\ trv1 \wedge lcompletedFromV\ sv1\ trv1 \wedge Van.lvalidFromS\ sv2$
 $trv2 \wedge$
 $lcompletedFromV\ sv2\ trv2 \wedge Van.lS\ trv1 = Opt.lS\ tr1 \wedge Van.lS\ trv2 = Opt.lS$
 $tr2 \wedge$
 $Van.lA\ trv1 = Van.lA\ trv2 \wedge Van.lO\ trv1 \neq Van.lO\ trv2$
apply(*rule exI[of - sv1]*)
apply(*rule exI[of - ltrv1 Δ ($\infty, \infty, s1, tr1, s2, tr2, Eq, sv1, sv2, Eq$)]*)
apply(*rule exI[of - sv2]*)
apply(*rule exI[of - ltrv2 Δ ($\infty, \infty, s1, tr1, s2, tr2, Eq, sv1, sv2, Eq$)]*)
apply(*intro conjI*)
subgoal by fact subgoal by fact subgoal by fact subgoal by fact subgoal by fact
subgoal apply(*rule lvalidFromS-ltrv1*) **using all by auto**
subgoal apply(*rule lcompletedFrom-ltrv1*) **using all by auto**
subgoal apply(*rule lvalidFromS-ltrv2*) **using all by auto**
subgoal apply(*rule lcompletedFrom-ltrv2*) **using all by auto**
subgoal apply(*rule lS-ltrv1-ltr1*) **using all by auto**
subgoal apply(*rule lS-ltrv2-ltr2*) **using all by auto**
subgoal apply(*rule lA-ltrv1-ltrv2*) **using all by auto**
subgoal using O34 apply- apply(*erule contrapos- nn*)
apply(*rule lO-ltrv1-ltrv2*) **using all by auto** .
qed

4.4 Compositional unwinding

We allow networks of unwinding relations that unwind into each other, which offer a compositional alternative to monolithic unwinding.

definition *unwindIntoCond* ::

$(enat \Rightarrow enat \Rightarrow enat \Rightarrow 'stateO \Rightarrow 'stateO \Rightarrow status \Rightarrow 'stateV \Rightarrow 'stateV \Rightarrow$
 $status \Rightarrow bool) \Rightarrow$
 $(enat \Rightarrow enat \Rightarrow enat \Rightarrow 'stateO \Rightarrow 'stateO \Rightarrow status \Rightarrow 'stateV \Rightarrow 'stateV \Rightarrow$
 $status \Rightarrow bool)$
 $\Rightarrow bool$

where

$unwindIntoCond\ \Delta\ \Delta' \equiv \forall w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$
 $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \longrightarrow$
 $(finalO\ s1 \longleftrightarrow finalO\ s2) \wedge (finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO$
 $s2)$
 \wedge
 $(statA = Eq \longrightarrow (isIntO\ s1 \longleftrightarrow isIntO\ s2))$
 \wedge
 $((\exists v < w. proact\ \Delta'\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO)$
 \vee
 $match\ \Delta'\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO)$

theorem *distrib-unwind-lrsecure*:
assumes $m: 0 < m$ **and** $\text{next}: \bigwedge i. i < (m::\text{nat}) \implies \text{next } i \subseteq \{0..<m\}$
and $\text{init}: \text{initCond } (\Delta s \ 0)$
and $\text{step}: \bigwedge i. i < m \implies$
 $\text{unwindIntoCond } (\Delta s \ i) (\lambda w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}.$
 $\exists j \in \text{next } i. \Delta s \ j \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO})$
shows *lrsecure*
proof –
define Δ **where** $D: \Delta \equiv \lambda w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}. \exists i < m. \Delta s \ i \ w$
 $w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}$
have $i: \text{initCond } \Delta$ **using** $\text{init } m$ **unfolding** *initCond-def* **by** *meson*
have $c: \text{unwindCond } \Delta$ **unfolding** *unwindCond-def* **apply**(*intro allI impI allI*)
apply(*subst (asm) D*) **apply** (*elim exE conjE*)
subgoal for $w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO} \ i$
apply(*frule step*) **unfolding** *unwindIntoCond-def*
apply(*erule allE[of - w]*) **apply**(*erule allE[of - w1]*) **apply**(*erule allE[of -*
 $w2]$)
apply(*erule allE[of - s1]*) **apply**(*erule allE[of - s2]*) **apply**(*erule allE[of -*
 $\text{statA}]$)
apply(*erule allE[of - sv1]*) **apply**(*erule allE[of - sv2]*) **apply**(*erule allE[of -*
 $\text{statO}]$)
apply *simp* **apply**(*elim conjE*)
apply(*erule disjE*)
subgoal **apply**(*rule disjI1*)
subgoal **apply**(*elim exE conjE*) **subgoal for** v
apply(*rule exI[of - v], simp*)
 $\Delta s \ j \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}. \exists j \in \text{next } i.$
 $\Delta s \ j \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO})$
subgoal **unfolding** *le-fun-def D* **by** *simp (meson atLeastLessThan-iff next*
 $\text{subsetD})$
subgoal
subgoal **apply**(*rule disjI2*)
 $\Delta s \ j \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}. \exists j \in \text{next } i.$
 $\Delta s \ j \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO})$
subgoal **unfolding** *le-fun-def D* **by** *simp (meson atLeastLessThan-iff next*
 $\text{subsetD})$
subgoal
show *?thesis using unwind-lrsecure[OF i c]* .
qed

lemma *unwindIntoCond-simpleI*:

assumes

$\bigwedge w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}.$

$\text{reachO } s1 \implies \text{reachO } s2 \implies \text{reachV } sv1 \implies \text{reachV } sv2 \implies$

$\Delta \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ \text{sv1} \ \text{sv2} \ \text{statO}$

\implies

$(\text{finalO } s1 \longleftrightarrow \text{finalO } s2) \wedge (\text{finalV } sv1 \longleftrightarrow \text{finalO } s1) \wedge (\text{finalV } sv2 \longleftrightarrow \text{finalO } s2)$

and

$\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $statA = Eq$

\implies

$isIntO s1 \longleftrightarrow isIntO s2$
 $\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$

\implies

$match \Delta' w1 w2 s1 s2 statA sv1 sv2 statO$

shows $unwindIntoCond \Delta \Delta'$

using *assms unfolding unwindIntoCond-def by auto*

lemma *unwindIntoCond-simpleI2:*

assumes

$\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$

\implies

$(finalO s1 \longleftrightarrow finalO s2) \wedge (finalV sv1 \longleftrightarrow finalO s1) \wedge (finalV sv2 \longleftrightarrow finalO s2)$

and

$\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $statA = Eq$

\implies

$isIntO s1 \longleftrightarrow isIntO s2$

and

$\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$

\implies

$(\exists v < w. proact \Delta' v w1 w2 s1 s2 statA sv1 sv2 statO)$

shows $unwindIntoCond \Delta \Delta'$

using *assms unfolding unwindIntoCond-def by auto*

lemma *unwindIntoCond-simpleIB:*

assumes

$\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$

\implies

$(finalO s1 \longleftrightarrow finalO s2) \wedge (finalV sv1 \longleftrightarrow finalO s1) \wedge (finalV sv2 \longleftrightarrow finalO s2)$

and

$\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$

$reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$
 $statA = Eq$
 \implies
 $isIntO\ s1 \longleftrightarrow isIntO\ s2$
and
 $\bigwedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
 \implies
 $(\exists v < w. proact\ \Delta'\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO) \vee match\ \Delta'\ w1\ w2\ s1\ s2$
 $statA\ sv1\ sv2\ statO$
shows *unwindIntoCond* $\Delta\ \Delta'$
using *assms unfolding unwindIntoCond-def by auto*

definition *oor where*

$oor\ \Delta\ \Delta_2 \equiv \lambda w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \vee \Delta_2\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

lemma *oorI1:*

$\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies oor\ \Delta\ \Delta_2\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2$
 $statO$

unfolding *oor-def by simp*

lemma *oorI2:*

$\Delta_2\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies oor\ \Delta\ \Delta_2\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2$
 $statO$

unfolding *oor-def by simp*

definition *oor3 where*

$oor3\ \Delta\ \Delta_2\ \Delta_3 \equiv \lambda w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \vee \Delta_2\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \vee$
 $\Delta_3\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

lemma *oor3I1:*

$\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies oor3\ \Delta\ \Delta_2\ \Delta_3\ w\ w1\ w2\ s1\ s2\ statA\ sv1$
 $sv2\ statO$

unfolding *oor3-def by simp*

lemma *oor3I2:*

$\Delta_2\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies oor3\ \Delta\ \Delta_2\ \Delta_3\ w\ w1\ w2\ s1\ s2\ statA$
 $sv1\ sv2\ statO$

unfolding *oor3-def by simp*

lemma *oor3I3:*

$\Delta_3\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies oor3\ \Delta\ \Delta_2\ \Delta_3\ w\ w1\ w2\ s1\ s2\ statA$
 $sv1\ sv2\ statO$

unfolding *oor3-def* by *simp*

definition *oor4* where

$oor4 \Delta \Delta_2 \Delta_3 \Delta_4 \equiv \lambda w w1 w2 s1 s2 statA sv1 sv2 statO.$

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO$

lemma *oor4I1*:

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w w1 w2 s1 s2 statA$
 $sv1 sv2 statO$

unfolding *oor4-def* by *simp*

lemma *oor4I2*:

$\Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w w1 w2 s1 s2 statA$
 $sv1 sv2 statO$

unfolding *oor4-def* by *simp*

lemma *oor4I3*:

$\Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w w1 w2 s1 s2 statA$
 $sv1 sv2 statO$

unfolding *oor4-def* by *simp*

lemma *oor4I4*:

$\Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w w1 w2 s1 s2 statA$
 $sv1 sv2 statO$

unfolding *oor4-def* by *simp*

definition *oor5* where

$oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \equiv \lambda w w1 w2 s1 s2 statA sv1 sv2 statO.$

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO$
 \vee
 $\Delta_5 w w1 w2 s1 s2 statA sv1 sv2 statO$

lemma *oor5I1*:

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w w1 w2 s1 s2$
 $statA sv1 sv2 statO$

unfolding *oor5-def* by *simp*

lemma *oor5I2*:

$\Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w w1 w2 s1 s2$
 $statA sv1 sv2 statO$

unfolding *oor5-def* by *simp*

lemma *oor5I3*:

$\Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w w1 w2 s1 s2$
 $statA sv1 sv2 statO$

unfolding *oor5-def* by *simp*

lemma *oor5I4*:

$\Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w w1 w2 s1 s2$
 $statA sv1 sv2 statO$

unfolding *oor5-def* **by** *simp*

lemma *oor5I5*:

$\Delta_5 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w w1 w2 s1 s2$
 $statA sv1 sv2 statO$

unfolding *oor5-def* **by** *simp*

lemma *isIntO-match1*: $isIntO s1 \implies match1 \Delta w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *match1-def* **by** *auto*

lemma *isIntO-match2*: $isIntO s2 \implies match2 \Delta w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *match2-def* **by** *auto*

lemma *isIntO-match*:

assumes $\langle isIntO s1 \rangle$ **and** $\langle isIntO s2 \rangle$

and $\langle match12 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \rangle$

shows $\langle match \Delta w1 w2 s1 s2 statA sv1 sv2 statO \rangle$

unfolding *match-def* **apply** (*intro conjI*)

subgoal

using *assms(1)* **by** (*rule isIntO-match1*)

subgoal

using *assms(2)* **by** (*rule isIntO-match2*)

subgoal

using *assms(3)* **by** *assumption*

.

lemma *match1-1-oorI1*:

$match1-1 \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \implies$

$match1-1 (oor \Delta \Delta_2) w1 w2 s1 s1' s2 statA sv1 sv2 statO$

apply(*rule match1-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match1-1-oorI2*:

$match1-1 \Delta_2 w1 w2 s1 s1' s2 statA sv1 sv2 statO \implies$

$match1-1 (oor \Delta \Delta_2) w1 w2 s1 s1' s2 statA sv1 sv2 statO$

apply(*rule match1-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match1-oorI1*:

$match1 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \implies$

$match1 (oor \Delta \Delta_2) w1 w2 s1 s2 statA sv1 sv2 statO$

apply(*rule match1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match1-oorI2*:
 $match1 \Delta_2 w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $match1 (oor \Delta \Delta_2) w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *match1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match2-1-oorI1*:
 $match2-1 \Delta w1 w2 s1 s2 s2' statA sv1 sv2 statO \implies$
 $match2-1 (oor \Delta \Delta_2) w1 w2 s1 s2 s2' statA sv1 sv2 statO$
apply(rule *match2-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match2-1-oorI2*:
 $match2-1 \Delta_2 w1 w2 s1 s2 s2' statA sv1 sv2 statO \implies$
 $match2-1 (oor \Delta \Delta_2) w1 w2 s1 s2 s2' statA sv1 sv2 statO$
apply(rule *match2-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match2-oorI1*:
 $match2 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $match2 (oor \Delta \Delta_2) w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *match2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match2-oorI2*:
 $match2 \Delta_2 w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $match2 (oor \Delta \Delta_2) w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *match2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-oorI1*:
 $match12 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $match12 (oor \Delta \Delta_2) w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *match12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-oorI2*:
 $match12 \Delta_2 w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $match12 (oor \Delta \Delta_2) w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *match12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-1-oorI1*:
 $match12-1 \Delta w1 w2 s1' s2' statA' sv1 sv2 statO \implies$
 $match12-1 (oor \Delta \Delta_2) w1 w2 s1' s2' statA' sv1 sv2 statO$
apply(rule *match12-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-1-oorI2*:
 $match12-1 \Delta_2 w1 w2 s1' s2' statA' sv1 sv2 statO \implies$
 $match12-1 (oor \Delta \Delta_2) w1 w2 s1' s2' statA' sv1 sv2 statO$
apply(rule *match12-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-2-oorI1*:
 $match12-2 \Delta w1 w2 s2 s2' statA' sv1 sv2 statO \implies$
 $match12-2 (oor \Delta \Delta_2) w1 w2 s2 s2' statA' sv1 sv2 statO$
apply(rule *match12-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-2-oorI2*:
 $match12-2 \Delta_2 w1 w2 s2 s2' statA' sv1 sv2 statO \implies$
 $match12-2 (oor \Delta \Delta_2) w1 w2 s2 s2' statA' sv1 sv2 statO$
apply(rule *match12-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-12-oorI1*:
 $match12-12 \Delta w1 w2 s1' s2' statA' sv1 sv2 statO \implies$
 $match12-12 (oor \Delta \Delta_2) w1 w2 s1' s2' statA' sv1 sv2 statO$
apply(rule *match12-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-12-oorI2*:
 $match12-12 \Delta_2 w1 w2 s1' s2' statA' sv1 sv2 statO \implies$
 $match12-12 (oor \Delta \Delta_2) w1 w2 s1' s2' statA' sv1 sv2 statO$
apply(rule *match12-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match-oorI1*:
 $match \Delta w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $match (oor \Delta \Delta_2) w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *match-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match-oorI2*:
 $match \Delta_2 w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $match (oor \Delta \Delta_2) w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *match-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *proact-oorI1*:
 $proact \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $proact (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *proact-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *proact-oorI2*:
 $proact \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $proact (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *proact-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-1-oorI1*:
 $move-1 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-1 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-1-oorI2*:
 $move-1 \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-1 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-2-oorI1*:
 $move-2 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-2 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-2-oorI2*:
 $move-2 \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-2 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-12-oorI1*:
 $move-12 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-12 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-12-oorI2*:
 $move-12 \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-12 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

end

context *Relative-Security-Determ*
begin

lemma *match1-1-defD*: $match1-1 \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv1 \wedge \Delta \infty w1 w2 s1' s2 statA (nextO sv1) sv2 statO$
unfolding *match1-1-def validTrans-iff-next* **by** *simp*

lemma *match1-12-defD*: $match1-12 \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv1 \wedge \neg finalV sv2 \wedge$
 $\Delta \infty w1 w2 s1' s2 statA (nextO sv1) (nextO sv2) (sstatO' statO sv1 sv2)$
unfolding *match1-12-def validTrans-iff-next* **by** *simp*

lemmas *match1-defsD* = *match1-def match1-1-defD match1-12-defD*

lemma *match2-1-defD*: $match2-1 \Delta w1 w2 s1 s2 s2' statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv2 \wedge \Delta \infty w1 w2 s1 s2' statA sv1 (nextO sv2) statO$
unfolding *match2-1-def validTrans-iff-next* **by** *simp*

lemma *match2-12-defD*: $match2-12 \Delta w1 w2 s1 s2 s2' statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv1 \wedge \neg finalV sv2 \wedge \Delta \infty w1 w2 s1 s2' statA (nextO sv1) (nextO sv2)$
 $(sstatO' statO sv1 sv2)$
unfolding *match2-12-def validTrans-iff-next* **by** *simp*

lemmas *match2-defsD* = *match2-def match2-1-defD match2-12-defD*

lemma *match12-1-defD*: $match12-1 \Delta w1 w2 s1' s2' statA' sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv1 \wedge \Delta \infty w1 w2 s1' s2' statA' (nextO sv1) sv2 statO$
unfolding *match12-1-def validTrans-iff-next* **by** *simp*

lemma *match12-2-defD*: $match12-2 \Delta w1 w2 s1' s2' statA' sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv2 \wedge \Delta \infty w1 w2 s1' s2' statA' sv1 (nextO sv2) statO$
unfolding *match12-2-def validTrans-iff-next* **by** *simp*

lemma *match12-12-defD*: $match12-12 \Delta w1 w2 s1' s2' statA' sv1 sv2 statO \longleftrightarrow$

$(let statO' = sstatO' statO sv1 sv2 in$
 $\neg finalV sv1 \wedge \neg finalV sv2 \wedge$
 $(statA' = Diff \longrightarrow statO' = Diff) \wedge$
 $\Delta \infty w1 w2 s1' s2' statA' (nextO sv1) (nextO sv2) statO')$

unfolding *match12-12-def validTrans-iff-next* **by** *simp*

lemmas *match12-defsD* = *match12-def match12-1-defD match12-2-defD match12-12-defD*

lemmas *match-deep-defsD* = *match1-defsD match2-defsD match12-defsD*

lemma *move-1-defD*: $move-1 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv1 \wedge \Delta w w1 w2 s1 s2 statA (nextO sv1) sv2 statO$
unfolding *move-1-def validTrans-iff-next* **by** *simp*

lemma *move-2-defD*: $move-2 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv2 \wedge \Delta w w1 w2 s1 s2 statA sv1 (nextO sv2) statO$
unfolding *move-2-def validTrans-iff-next* **by** *simp*

lemma *move-12-defD*: $move-12 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \longleftrightarrow$
 $(let statO' = sstatO' statO sv1 sv2 in$
 $\neg finalV sv1 \wedge \neg finalV sv2 \wedge$
 $\Delta w w1 w2 s1 s2 statA (nextO sv1) (nextO sv2) statO')$
unfolding *move-12-def validTrans-iff-next* **by** *simp*

lemmas *proact-defsD* = *proact-def move-1-defD move-2-defD move-12-defD*

end

end

References

- [1] A. P. Brijesh Dongol, Matt Griffin and J. Wright. Relative security: Formally modeling and (dis)proving resilience against semantic optimization vulnerabilities. In *37th IEEE Computer Security Foundations Symposium, CSF 2024*. To appear.