Proof Terms for Term Rewriting

Christina Kirk (Kohl)

University of Innsbruck, Austria

April 13, 2025

Abstract

Proof terms are first-order terms that represent reductions in term rewriting. They were initially introduced in [6] and [5, Chapter 8] by van Oostrom and de Vrijer to study equivalences of reductions in left-linear rewrite systems. This entry formalizes proof terms for multi-steps in first-order term rewrite systems. We define simple proof terms (i.e., without a composition operator) and establish the correspondence to multi-steps: each proof term represents a multi-step with the same source and target, and every multi-step can be expressed as a proof term. The formalization moreover includes operations on proof terms, such as residuals, join, and deletion and a method for labeling proof term sources to identify overlaps between two proof terms.

This formalization is part of the *Isabelle Formalization of Rewriting* IsaFoR and is an essential component of several formalized confluence and commutation results involving multi-steps [2, 3, 4, 1].

Contents

1 Preliminaries **2**

1.1 Utilities for Lists 2

1.1.1 Lists of *option* 8

1.2 Results About Linear Terms 9

1.3 Results About Substitutions and Contexts 12

1.3.1 Utilities for *mk-subst* 17

1.4 Matching Terms 19

1.4.1 Matching of Linear Terms 21

2 Proof Terms **25**

2.1 Definitions 25

2.2 Frequently Used Locales/Contexts 27

2.3 Proof Term Predicates 28

2.4 'Normal' Terms vs. Proof Terms 32

2.5 Substitutions . . . 33
2.6 Contexts . . . 37
2.7 Source and Target . . . 39
2.8 Additional Results . . . 49
2.9 Proof Terms Represent Multi-Steps . . . 51

3 Operations on Proof Terms **53**
3.1 Residuals . . . 57
3.2 Join . . . 73
3.2.1 N-Fold Join . . . 89
3.3 Deletion . . . 125
3.4 Computations With Single Redexes . . . 128

4 Orthogonal Proof Terms **138**

5 Labels and Overlaps **150**
5.1 Labeled Proof Terms . . . 150
5.2 Measuring Overlap . . . 199
5.3 Collecting Overlapping Positions . . . 208

6 Redex Patterns **212**

1 Preliminaries

theory *Proof-Term-Utils*
imports
First-Order-Terms.Matching
First-Order-Rewriting.Term-Impl
begin

1.1 Utilities for Lists

lemma *obtain-list-with-property*:
assumes $\forall x \in set\ xs.\ \exists a.\ P\ a\ x$
shows $\exists as.\ length\ as = length\ xs \wedge (\forall i < length\ xs.\ P\ (as!i)\ (xs!i))$
using *assms* **proof**(*induct xs*)
case (*Cons a xs*)
then show *?case*
by (*metis length-map nth-map nth-mem*)
qed *simp*

lemma *card-Union-Sum*:
assumes *is-partition* (*map f* [*0*..<*length xs*])
and $\forall i < length\ xs.\ finite\ (f\ i)$
shows $card\ (\bigcup i{<}length\ xs.\ f\ i) = (\sum i{<}length\ xs.\ card\ (f\ i))$
proof−
from *assms*(*1*) **have** $disj{:}pairwise\ (\lambda s\ t.\ disjnt\ (f\ s)\ (f\ t))\ \{..{<}length\ xs\}$

unfolding *pairwise-def is-partition-alt is-partition-alt-def disjnt-def* **by** *simp*
then have *pairwise disjnt* $(f$ ‘ $\{..<$*length xs*$\})$
by (*metis* (*mono-tags*, *lifting*) *pairwiseD pairwise-imageI*)
then have *card* $(\bigcup i<$*length xs*. f $i)$ = *sum card* $(f$ ‘ $\{..<$*length xs*$\})$
using *assms*(*2*) *card-Union-disjoint* **by** (*metis* (*mono-tags*, *lifting*) *imageE lessThan-iff*)
with *disj* **show** *?thesis*
using *sum-card-image* **by** (*metis finite-lessThan*)
qed

lemma *sum-sum-concat*: $(\sum i<$*length xs*. $\sum x\leftarrow f$ $(xs!i)$. g $x)$ = $(\sum x\leftarrow$*concat* (*map* f *xs*). g $x)$
proof(*induct xs*)
case (*Cons a xs*)
then show *?case* **unfolding** *list.map concat.simps map-append sum-list-append*
by (*metis* (*mono-tags*, *lifting*) *length-nth-simps*(*2*) *nth-Cons-0 nth-Cons-Suc sum.cong sum.lessThan-Suc-shift*)
qed *simp*

lemma *concat-map2-zip*:
assumes *length xs* = *length ys*
and $\forall i <$ *length xs*. *length* $(xs!i)$ = *length* $(ys!i)$
shows *concat* (*map2 zip xs ys*) = *zip* (*concat xs*) (*concat ys*)
using *assms* **proof**(*induct xs arbitrary*:*ys rule*:*rev-induct*)
case (*snoc x xs*)
from *snoc*(*2*) **obtain** y ys' **where** y:$ys = ys'$@$[y]$
by (*metis append-is-Nil-conv length-0-conv neq-Nil-conv rev-exhaust*)
moreover with *snoc*(*2*) **have** l:*length xs* = *length* ys' **by** *simp*
moreover with *snoc*(*3*) **have** l':$\forall i <$ *length xs*. *length* $(xs!i)$ = *length* $(ys'!i)$
unfolding y **by** (*metis* (*no-types*, *lifting*) *Ex-less-Suc add-Suc-right append.right-neutral append-Cons-nth-left length-Cons length-append order-less-trans*)
ultimately have *IH*:*concat* (*map2 zip xs* ys') = *zip* (*concat xs*) (*concat* ys')
using *snoc*(*1*) **by** *presburger*
have $*$:*concat* (*map2 zip* (*xs* @ $[x]$) *ys*) = *concat* (*map2 zip xs* ys') @ (*zip x y*)
unfolding y *zip-append*[*OF* l] **by** *simp*
have *length* (*concat xs*) = *length* (*concat* ys')
using l l' *eq-length-concat-nth* **by** *blast*
then show *?case*
unfolding $*$ *IH* **unfolding** y *concat-append* **using** *zip-append* **by** *simp*
qed *simp*

lemma *sum-list-less*:
assumes *less*:$i < j$
and $i'j'$:$i' <$ *length xs* $j' <$ *length xs*
and j'':$j'' <$ *length* $(xs!j')$
and *sums*:i = *sum-list* (*map length* (*take* i' *xs*)) + i'' j = *sum-list* (*map length* (*take* j' *xs*)) + j''
shows $i' \leq j'$
proof(*rule ccontr*)

assume $*$:$\neg$ $i' \leq j'$
then have *subsums*:*sum-list* (*map length* (*take* i' *xs*)) = *sum-list* (*map length* (*take* j' *xs*)) + *sum-list* (*map length* (*take* (i'−j') (*drop* j' *xs*)))
by (*metis le-add-diff-inverse map-append nat-le-linear sum-list-append take-add*)

from $*$ **have** *take* ($i' - j'$) (*drop* j' *xs*) = *xs*!j' # (*take* (i' − (*Suc* j')) (*drop* (*Suc* j') *xs*))
using $i'j'$ **by** (*metis Cons-nth-drop-Suc Suc-diff-Suc linorder-le-less-linear take-Suc-Cons*)

with j'' **have** j'' < *sum-list* (*map length* (*take* (i'−j') (*drop* j' *xs*))) **by** *simp*
then show *False*
using *sums subsums less* **by** *linarith*
qed

lemma *zip-symm*: (x, y) ∈ *set* (*zip xs ys*) ⟹ (y, x) ∈ *set* (*zip ys xs*)
by (*induct xs ys rule*:*list-induct2*′) *auto*

lemma *sum-list-elem*:
($\sum$ x←[y]. f x) = f y
by *simp*

lemma *sum-list-zero*:
assumes ∀ i < *length xs*. f (*xs*!i) = *0*
shows ($\sum$ x←*xs*. f x) = *0*
by (*metis assms map-eq-conv*′ *monoid-add-class.sum-list-0*)

lemma *distinct-is-partition*:
assumes *distinct* (*concat ts*)
shows *is-partition* (*map set ts*)
using *assms* **proof**(*induct ts*)
case *Nil*
then show *?case*
using *is-partition-Nil* **by** *auto*
next
case (*Cons t ts*)
{**fix** i j **assume** j:j < *length* (t#*ts*) **and** *ij*:$i < j$
have (*map set* (t#*ts*))!i ∩ (*map set* (t#*ts*))!j = {} **proof**(*cases* i)
case *0*
show *?thesis* **using** *Cons*(*2*) **unfolding** *0*
using *ij j* **by** *force*
next
case (*Suc* n)
from *Cons* **have** *is-partition* (*map set ts*) **by** *simp*
then show *?thesis*
unfolding *Suc is-partition-def* **using** *j ij* **using** *Suc Suc-less-eq2* **by** *fastforce*
qed
}
then show *?case* **unfolding** *is-partition-def* **by** *simp*
qed

lemma *filter-ex-index*:
  assumes $x = filter\ f\ xs\ !\ i\ i < length\ (filter\ f\ xs)$
  shows $\exists j.\ j < length\ xs \wedge x = xs!j$
  using *assms* **proof**(*induct xs arbitrary*:*i*)
  case (*Cons y ys*)
  show *?case* **proof**(*cases f y*)
    case *True*
    then have *filter*:$filter\ f\ (y\#ys) = y\#(filter\ f\ ys)$ **by** *simp*
    show *?thesis* **proof**(*cases i*)
      case *0*
      from *Cons*(*2*) **show** *?thesis* **unfolding** *filter 0* **by** *auto*
    next
      case (*Suc n*)
      from *Cons*(*2*) **have** $x = filter\ f\ ys\ !\ n$
        unfolding *Suc filter* **by** *simp*
      moreover from *Cons*(*3*) **have** $n < length\ (filter\ f\ ys)$
        unfolding *Suc filter* **by** *simp*
      ultimately obtain *j* **where** $j<length\ ys$ **and** $x = ys\ !\ j$
        using *Cons*(*1*) **by** *blast*
      then show *?thesis* **by** *auto*
    qed
  next
    case *False*
    then have *filter*:$filter\ f\ (y\#ys) = filter\ f\ ys$ **by** *simp*
    from *Cons* **obtain** *j* **where** $j<length\ ys$ **and** $x = ys\ !\ j$
      unfolding *filter* **by** *blast*
    then show *?thesis* **by** *auto*
  qed
qed *simp*

lemma *filter-index-neq'*:
  assumes $i < j\ j < length\ (filter\ f\ xs)$
  shows $\exists\ i'\ j'.\ i' < length\ xs \wedge j' < length\ xs \wedge i' < j' \wedge xs\ !\ i' = (filter\ f\ xs)\ !\ i \wedge xs\ !\ j' = (filter\ f\ xs)\ !\ j$
  using *assms* **proof**(*induct xs arbitrary*: *i j*)
  case (*Cons x xs*)
  then show *?case* **proof**(*cases f x*)
    case *True*
    show *?thesis* **proof**(*cases i*)
      case *0*
      then have *i0*:$filter\ f\ (x\#xs)\ !\ i = (x\#xs)\ !\ 0$
        using ‹*f x*› **by** *simp*
      from *Cons*(*2*) **obtain** j' **where** $j = Suc\ j'$
        unfolding *0* **using** *gr0-implies-Suc* **by** *blast*
      with *Cons*(*3*) **have** $j' < length\ (filter\ f\ xs)$
        unfolding *filter.simps* **using** ‹*f x*› **by** *simp*
      then obtain j'' **where** j'':$j'' < length\ xs\ filter\ f\ xs\ !\ j' = xs\ !\ j''$
        by (*meson filter-ex-index*)

```
      then have filter f (x#xs) ! j = (x#xs) ! (Suc j'')
        using ‹f x› ‹j = Suc j'› by simp
      with i0 j''(1) show ?thesis
        by (metis length-nth-simps(2) not-less-eq zero-less-Suc)
    next
      case (Suc i')
      from Cons(2) obtain j' where j:j = Suc j'
        unfolding Suc using Suc-lessE by auto
      from Cons(1)[of i' j'] Cons(2,3) obtain i'' j'' where i'' < length xs j'' <
length xs i'' < j'' xs ! i'' = filter f xs ! i' xs ! j'' = filter f xs ! j'
        using Suc True j by auto
      then show ?thesis
        by (smt (verit) Suc Suc-less-eq True filter.simps(2) j length-nth-simps(2)
nth-Cons-Suc)
    qed
  next
    case False
    then have filter f (x#xs) = filter f xs by simp
    with Cons show ?thesis
      by (metis Suc-less-eq length-nth-simps(2) nth-Cons-Suc)
  qed
qed simp

lemma filter-index-neq:
  assumes i ≠ j i < length (filter f xs) j < length (filter f xs)
  shows ∃ i' j'. i' < length xs ∧ j' < length xs ∧ i' ≠ j' ∧ xs ! i' = (filter f xs) !
i ∧ xs ! j' = (filter f xs) ! j
using assms filter-index-neq' proof(cases i < j)
  case False
  then have *:j < i using assms(1) by simp
  then show ?thesis using filter-index-neq'[OF * assms(2)] by blast
qed blast

lemma nth-drop-equal:
  assumes length xs = length ys
    and ∀ j < length xs. j ≥ i ⟶ xs!j = ys!j
  shows drop i xs = drop i ys
using assms proof (induct i arbitrary: xs ys)
  case 0
  then show ?case
    using nth-equalityI by blast
next
  case (Suc i)
  then show ?case proof(cases Suc i < length xs)
    case True
    then obtain x xs' where x:xs = x # xs'
      by (metis Suc-length-conv Suc-lessE)
    with Suc(2) obtain y ys' where y:ys = y # ys'
      by (metis length-greater-0-conv nth-drop-0)
```

```
    from Suc(1)[of xs' ys'] have drop i xs' = drop i ys'
      using Suc(2,3) unfolding x y
    by (metis Suc-le-mono length-nth-simps(2) linorder-not-le nat.inject nth-Cons-Suc)
    then show ?thesis unfolding x y by simp
  qed simp
qed

lemma union-take-drop-list:
  assumes i < length xs
  shows (set (take i xs)) ∪ (set (drop (Suc i) xs)) = {xs!j | j. j < length xs ∧ j ≠
i}
proof−
  from assms have i:i ≤ length xs by simp
  have set1:set (take i xs) = {xs ! j |j. j < i}
    using nth-image[OF i] unfolding image-def by fastforce
  from assms have i:Suc i ≤ length xs by simp
  have set2:set (drop (Suc i) xs) = {xs !j |j. i < j ∧ j < length xs} proof
    {fix x assume x ∈ set (drop (Suc i) xs)
      then have x ∈ {xs !j |j. i < j ∧ j < length xs}
        unfolding set-conv-nth nth-drop[OF i] length-drop by auto
    }
    then show set (drop (Suc i) xs) ⊆ {xs ! j |j. i < j ∧ j < length xs} by auto
    {fix x assume x ∈ {xs !j |j. i < j ∧ j < length xs}
      then have x ∈ set (drop (Suc i) xs)
        unfolding set-conv-nth nth-drop[OF i] length-drop
      by (smt (verit, best) Suc-leI add-diff-inverse-nat i mem-Collect-eq nat-add-left-cancel-less
not-less-eq-eq)
    }
    then show {xs ! j |j. i < j ∧ j < length xs} ⊆ set (drop (Suc i) xs) by auto
  qed
  {fix x assume x ∈ set (take i xs) ∪ set (drop (Suc i) xs)
    then consider x ∈ set (take i xs) | x ∈ set (drop (Suc i) xs) by blast
    then have x ∈ {xs ! j |j. j < length xs ∧ j ≠ i} proof(cases)
      case 1
      with set1 show ?thesis using in-set-idx by fastforce
    next
      case 2
      with set2 show ?thesis using in-set-idx by fastforce
    qed
  }moreover
  {fix x assume x ∈ {xs ! j |j. j < length xs ∧ j ≠ i}
    then obtain j where x = xs!j and j:j < length xs j ≠ i
      by blast
    then have x ∈ set (take i xs) ∪ set (drop (Suc i) xs)
      using set1 set2 using nat-neq-iff by auto
  }
  ultimately show ?thesis by auto
qed
```

lemma *list-tl-eq*:
assumes *length xs* = *length ys xs* ≠ []
and ∀ *i* < *length xs*. *i* > *0* ⟶ *xs*!*i* = *ys*!*i*
shows *tl xs* = *tl ys*
by (*metis Suc-le-lessD assms*(*1*) *assms*(*3*) *length-greater-0-conv list.sel*(*3*) *nth-drop-0 nth-drop-equal*)

1.1.1 Lists of *option*

lemma *length-those*:
assumes *those xs* = *Some ys*
shows *length xs* = *length ys*
using *assms* **proof**(*induction xs arbitrary*:*ys*)
case *Nil*
then show *?case* **by** *simp*
next
case (*Cons a xs*)
from *Cons*(*2*) **obtain** *ys′* **where** *ys′*:*those xs* = *Some ys′*
by (*smt not-None-eq option.case-eq-if option.simps*(*8*) *those.simps*(*2*))
from *Cons*(*2*) **obtain** *y* **where** *y*:*Some y* = *a*
by (*metis option.case-eq-if option.exhaust-sel option.simps*(*3*) *those.simps*(*2*))
from *y ys′* **have** *those* (*Cons a xs*) = *Some* (*Cons y ys′*)
by *auto*
then show *?case* **using** *Cons ys′*
by *auto*
qed

lemma *those-not-none-x*: *those xs* = *Some ys* ⟹ *x* ∈ *set xs* ⟹ *x* ≠ *None*
proof (*induction xs arbitrary*: *x ys*)
case (*Cons a xs*)
from *Cons*(*2*) **have** *a* ≠ *None* **using** *option.simps*(*4*) **by** *fastforce*
from *this Cons*(*2*) **have** *those xs* ≠ *None* **by** *auto*
then show *?case* **using** *Cons*(*1*,*3*) ‹*a* ≠ *None*› **by** *auto*
qed (*simp*)

lemma *those-not-none-xs*: *list-all* (λ*x*. *x* ≠ *None*) *xs* ⟹ *those xs* ≠ *None*
by (*induction xs*) *auto*

lemma *those-some*:
assumes *length xs* = *length ys* ∀ *i* < *length xs*. *xs*!*i* = *Some* (*ys*!*i*)
shows *those xs* = *Some ys*
using *assms* **by** (*induct rule*:*list-induct2*) (*simp*, *force*)

lemma *those-some2*:
assumes *those xs* = *Some ys*
shows ∀ *i* < *length xs*. *xs*!*i* = *Some* (*ys*!*i*)
proof−
from *assms* **have** *length xs* = *length ys* **by** (*simp add*: *length-those*)
then show *?thesis* **using** *assms* **proof**(*induction xs ys rule*:*list-induct2*)

case (*Cons x xs y ys*)
from *Cons*(*3*) **have** $x \neq$ *None* **by** (*metis list.set-intros*(*1*) *those-not-none-x*)
with *Cons*(*3*) **have** $*$:$x =$ *Some y* **by** *force*
with *Cons*(*3*) **have** *those xs* = *Some ys* **by** *force*
with $*$ *Cons*(*2*) **show** *?case* **by** (*simp add*: *nth-Cons′*)
qed *simp*
qed

lemma *exists-some-list*:
assumes $\forall$ $i <$ *length xs*. ($\exists$ *y*. *xs*!*i* = *Some y*)
shows $\exists$ *ys*. ($\forall$ $i <$ *length xs*. *xs*!*i* = *Some* (*ys*!*i*)) $\wedge$ *length ys* = *length xs*
by (*metis assms length-map nth-map option.sel*)

1.2 Results About Linear Terms

lemma *linear-term-var-vars-term-list*:
assumes *linear-term t*
shows *vars-term-list t* = *vars-distinct t*
using *assms linear-term-distinct-vars*
by (*metis comp-apply distinct-rev remdups-id-iff-distinct rev-rev-ident*)

lemma *linear-term-unique-vars*:
assumes *linear-term s*
and $p \in$ *poss s* **and** *s*|-*p* = *Var x*
and $q \in$ *poss s* **and** *s*|-*q* = *Var x*
shows $p = q$
proof(*rule ccontr*)
assume $p \neq q$
with *assms*(*2*−) **obtain** *i j* **where** *ij*:$i <$ *length* (*var-poss-list s*) $j <$ *length* (*var-poss-list s*) $i \neq j$
var-poss-list s ! $i = p$ *var-poss-list s* ! $j = q$
by (*metis in-set-idx var-poss-iff var-poss-list-sound*)
with *assms*(*3*,*5*) **have** *vars-term-list s* ! *i* = *vars-term-list s* ! *j*
by (*metis length-var-poss-list term.inject*(*1*) *vars-term-list-var-poss-list*)
moreover from *assms*(*1*) **have** *distinct* (*vars-term-list s*)
by (*metis distinct-remdups distinct-rev linear-term-var-vars-term-list o-apply*)
ultimately show *False* **using** *ij*(*1*,*2*,*3*)
by (*metis distinct-Ex1 length-var-poss-list nth-mem*)
qed

lemma *linear-term-ctxt*:
assumes *linear-term t*
and $p \in$ *poss t*
shows *vars-ctxt* (*ctxt-of-pos-term p t*) $\cap$ *vars-term* (*t*|-*p*) = {}
using *assms* **proof**(*induct p arbitrary*:*t*)
case (*Cons i p*)
from *Cons*(*3*) **obtain** *f ts* **where** *t*:*t* = *Fun f ts* $i <$ *length ts* $p \in$ *poss* (*ts*!*i*)
using *args-poss* **by** *blast*
with *Cons*(*1*,*2*) **have** *IH*:*vars-ctxt* (*ctxt-of-pos-term p* (*ts*!*i*)) $\cap$ *vars-term* ((*ts*!*i*)

|- *p*) = {}
  by *simp*
 {**fix** *j* **assume** *j*:*j* < *length ts j* ≠ *i*
  with *Cons*(*2*) **have** *vars-term* (*ts*!*j*) ∩ *vars-term* (*ts*!*i* |-*p*) = {}
   unfolding *t* **using** *var-in-linear-args t*(*2*,*3*) **by** (*metis* (*no-types*, *opaque-lifting*) *Int-Un-distrib disjoint-iff sup-bot.neutr-eq-iff vars-ctxt-pos-term*)
 }
 then have ⋃ {*vars-term* (*ts* ! *j*) |*j*. *j* < *length ts* ∧ *j* ≠ *i*} ∩ *vars-term* (*ts*!*i* |-*p*) = {}
  by *blast*
 moreover have (⋃ (*vars-term* ' *set* (*take i ts*)) ∪ ⋃ (*vars-term* ' *set* (*drop* (*Suc* *i*) *ts*))) =
   ⋃ {*vars-term* (*ts* ! *j*) |*j*. *j* < *length ts* ∧ *j* ≠ *i*}
  unfolding *set-map*[*symmetric*] *take-map*[*symmetric*] *drop-map*[*symmetric*] *Union-Un-distrib*[*symmetric*]
  using *union-take-drop-list*[**where** *xs*=(*map vars-term ts*)] **unfolding** *length-map* **using** *t*(*2*) **by** *auto*
 ultimately show *?case* **unfolding** *t ctxt-of-pos-term.simps subt-at.simps* **using** *IH*
   by (*metis* (*no-types*, *lifting*) *bot-eq-sup-iff inf-sup-distrib2 vars-ctxt.simps*(*2*))

qed *simp*

lemma *linear-term-obtain-subst*:
 assumes *linear-term* (*Fun f ts*) **and** *l*:*length ts* = *length ss*
  and *substs*: ∀ *i*< *length ts*. (∃ σ. *ts*!*i* · σ = *ss*!*i*)
 shows ∃ σ. *Fun f ts* · σ = *Fun f ss*
 using *assms* **proof**(*induct ts arbitrary*: *ss*)
 case (*Cons t ts*)
 from *Cons*(*3*) **obtain** *s ss'* **where** *ss*:*ss* = *s*#*ss'*
  by (*metis length-Suc-conv*)
 from *Cons*(*2*) **have** *lin*:*linear-term* (*Fun f ts*)
  unfolding *linear-term.simps* **by** (*simp add*: *is-partition-Cons*)
 from *Cons*(*4*) **have** ∀ *i*<*length ts*. ∃ σ. *ts* ! *i* · σ = *ss'* ! *i*
  unfolding *ss* **by** (*metis length-nth-simps*(*2*) *not-less-eq nth-Cons-Suc*)
 then obtain σ **where** σ:*Fun f ts* · σ = *Fun f ss'*
  using *Cons*(*1*)[*OF lin*, *of ss'*] **using** *Cons.prems*(*2*) *ss* **by** *auto*
 from *Cons*(*4*) **obtain** σ*1* **where** σ*1*:*t* · σ*1* = *s*
  using *ss* **by** *auto*
 let *?σ*=λ*x*. *if x* ∈ *vars-term t then* σ*1 x else* σ *x*
 have *t*:*t* · *?σ* = *s*
  by (*simp add*: σ*1 term-subst-eq*)
 {**fix** *i* **assume** *i* < *length ts*
  then have *ts*!*i* · *?σ* = *ss'*!*i*
   by (*smt* (*verit*, *ccfv-SIG*) *Cons.prems*(*1*) *Cons.prems*(*2*) *Suc-inject Suc-leI* σ σ*1 eval-term.simps*(*2*) *le-imp-less-Suc length-nth-simps*(*2*) *map-nth-eq-conv nth-Cons-0 nth-Cons-Suc ss term.sel*(*4*) *term-subst-eq var-in-linear-args zero-less-Suc*)
 }
 with *t* **have** *Fun f* (*t*#*ts*) · *?σ* = *Fun f ss*
  using *Cons.prems*(*2*) *map-nth-eq-conv ss* **by** *auto*

then show *?case* **by** *blast*
qed *simp*

lemma *linear-ctxt-of-pos-term*:
assumes *linear-term t* **and** *lin-s*:*linear-term s* **and** *p*:$p \in$ *poss t*
and *vars-term t* $\cap$ *vars-term s* = {}
shows *linear-term* (*replace-at t p s*)
using *assms* **proof**(*induct t arbitrary*:*p*)
case (*Var x*)
with *p* **have** *p* = [] **by** *simp*
with *lin-s* **show** *?case* **by** *simp*
next
case (*Fun f ts*)
from *lin-s* **show** *?case* **proof**(*cases p*)
case (*Cons i p'*)
with *Fun*(*4*) **have** *i*:*i* < *length ts* **by** *simp*
with *Fun*(*4*) **have** *p'*:*p'* $\in$ *poss* (*ts*!*i*) **unfolding** *Cons* **by** *simp*
{**fix** *n* **assume** *n*:*n* < *length ts n* $\neq$ *i*
with *Fun*(*2*) **have** *vars-term* (*ts*!*n*) $\cap$ *vars-term* (*ts*!*i*) = {}
by (*metis disjoint-iff i var-in-linear-args*)
then have *vars-term* (*ts*!*n*) $\cap$ *vars-ctxt* (*ctxt-of-pos-term p'* (*ts*!*i*)) = {}
using *p' vars-ctxt-pos-term* **by** *fastforce*
moreover from *n Fun*(*5*) **have** *vars-term* (*ts*!*n*) $\cap$ *vars-term s* = {}
by (*meson disjoint-iff nth-mem term.set-intros*(*4*))
ultimately have *vars-term* (*ts*!*n*) $\cap$ *vars-term* ((*ctxt-of-pos-term p'* (*ts* ! *i*))⟨*s*⟩) = {}
unfolding *vars-term-ctxt-apply* **by** *blast*
}
with *Fun*(*2*) **have** *is-partition* (*map vars-term* (*take i ts* @ (*ctxt-of-pos-term p'* (*ts* ! *i*))⟨*s*⟩ # *drop* (*Suc i*) *ts*))
unfolding *linear-term.simps is-partition-def* **by** (*smt* (*z3*) *Int-commute append-Cons-nth-not-middle i id-take-nth-drop*
length-append length-map length-nth-simps(*2*) *linorder-neq-iff nth-append-take nth-map order.strict-implies-order order.strict-trans*)
moreover have *linear-term* ((*ctxt-of-pos-term p'* (*ts* ! *i*))⟨*s*⟩)
using *Fun p'* **by** (*meson disjoint-iff i linear-term.simps*(*2*) *nth-mem term.set-intros*(*4*))

ultimately show *?thesis*
using *Fun*(*2*) **unfolding** *Cons ctxt-of-pos-term.simps intp-actxt.simps linear-term.simps*
by (*metis Un-iff in-set-dropD in-set-takeD set-ConsD set-append*)
qed *simp*
qed

lemma *distinct-vars*:
assumes $\bigwedge p\ q\ x\ y.\ p \neq q \Longrightarrow p \in$ *poss t* $\Longrightarrow q \in$ *poss t* $\Longrightarrow$ *t*|-*p* = *Var x* $\Longrightarrow$ *t*|-*q* = *Var y* $\Longrightarrow x \neq y$
shows *distinct* (*vars-term-list t*)
proof−

{fix *i j* **assume** *ij*:*i* ≠ *j* **and** *i*:*i* < *length* (*vars-term-list t*) **and** *j*:*j* < *length* (*vars-term-list t*)
let *?p=var-poss-list t* ! *i* **and** *?q=var-poss-list t* ! *j*
let *?x=vars-term-list t* ! *i* **and** *?y=vars-term-list t* ! *j*
from *ij i j* **have** *pq*:*?p* ≠ *?q*
by (*simp add*: *distinct-var-poss-list length-var-poss-list nth-eq-iff-index-eq*)
have *p*:*?p* ∈ *poss t*
by (*metis i length-var-poss-list nth-mem var-poss-imp-poss var-poss-list-sound*)

have *q*:*?q* ∈ *poss t*
by (*metis j length-var-poss-list nth-mem var-poss-imp-poss var-poss-list-sound*)

have *?x* ≠ *?y*
using *assms*[*OF pq p q*] *i j* **by** (*simp add*: *vars-term-list-var-poss-list*)
}
then show *?thesis* **by** (*meson distinct-conv-nth*)
qed

lemma *distinct-vars-linear-term*:
assumes *distinct* (*vars-term-list t*)
shows *linear-term t*
using *assms* **proof**(*induct t*)
case (*Fun f ts*)
{fix *t* **assume** *t*:*t* ∈ *set ts*
with *Fun*(*2*) **have** *distinct* (*vars-term-list t*)
unfolding *vars-term-list.simps* **by** (*simp add*: *distinct-concat-iff*)
with *t Fun*(*1*) **have** *linear-term t*
by *auto*
}
moreover have *is-partition* (*map vars-term ts*)
using *Fun*(*2*) **unfolding** *vars-term-list.simps* **using** *distinct-is-partition set-vars-term-list*
by (*metis* (*mono-tags*, *lifting*) *length-map map-nth-eq-conv*)
ultimately show *?case* **by** *simp*
qed *simp*

lemma *distinct-vars-eq-linear*: *linear-term t* = *distinct* (*vars-term-list t*)
using *distinct-vars-linear-term linear-term-distinct-vars* **by** *blast*

1.3 Results About Substitutions and Contexts

lemma *ctxt-apply-term-subst*:
assumes *linear-term t* **and** *i* < *length* (*vars-term-list t*)
and *p* = (*var-poss-list t*)!*i*
shows (*ctxt-of-pos-term p* (*t* · σ))⟨*s*⟩ = *t* · σ((*vars-term-list t*)!*i* := *s*)
proof−
from *assms*(*2*,*3*) **have** *t*|-*p* = *Var* ((*vars-term-list t*)!*i*)
by (*metis vars-term-list-var-poss-list*)
with *assms* **show** *?thesis*
by (*smt* (*verit*, *ccfv-threshold*) *filter-cong fun-upd-other fun-upd-same length-var-poss-list*

linear-term-replace-in-subst nth-mem var-poss-imp-poss var-poss-list-sound)
qed

lemma *ctxt-subst-apply*:
  assumes $p \in$ *poss t* **and** $t|\text{-}p =$ *Var x*
    and *linear-term t*
  shows ((*ctxt-of-pos-term p t*) $\cdot_c$ σ)$\langle s \rangle$ = $t \cdot \sigma(x := s)$
  unfolding *ctxt-of-pos-term-subst*[*symmetric*, *OF assms*(*1*)]
  using *assms*
  by (*smt* (*verit*) *fun-upd-apply linear-term-replace-in-subst*)

lemma *ctxt-of-pos-term-hole-subst*:
  assumes *linear-term t*
    and $i <$ *length* (*var-poss-list t*) **and** $p =$ *var-poss-list t* ! *i*
    and $\forall x \in$ *vars-term t*. $x \neq$ *vars-term-list t* !*i* $\longrightarrow \sigma\ x = \tau\ x$
  shows *ctxt-of-pos-term p* ($t \cdot \sigma$) = *ctxt-of-pos-term p* ($t \cdot \tau$)
  using *assms* **proof**(*induct p arbitrary*: *t i*)
  case (*Cons j p*)
  from *Cons*(*3*,*4*) **have** *j#p* $\in$ *var-poss t*
    using *nth-mem* **by** *force*
  then obtain *f ts* **where** *ts*:$j <$ *length ts* $t =$ *Fun f ts* $p \in$ *var-poss* (*ts*!*j*)
    by (*metis args-poss subt-at.simps*(*2*) *var-poss-iff*)
  then obtain i' **where** i':$i' <$ *length* (*var-poss-list* (*ts*!*j*)) $p =$ *var-poss-list* (*ts*!*j*)!i'
    using *var-poss-list-sound* **by** (*metis in-set-conv-nth*)
  from *Cons*(*3*,*4*) **have** *Var* (*vars-term-list t* ! *i*) = $t|$-(*j#p*)
    by (*metis length-var-poss-list vars-term-list-var-poss-list*)
  also have ... = (*ts*!*j*)|-*p*
    unfolding *ts*(*2*) **by** *simp*
  also have ... = *Var* (*vars-term-list* (*ts*!*j*) ! i')
    using i' **by** (*simp add*: *length-var-poss-list vars-term-list-var-poss-list*)
  finally have $*$:*vars-term-list t* ! i = *vars-term-list* (*ts* ! *j*) ! i' **by** *simp*
  with *Cons*(*5*) **have** $\forall x \in$*vars-term* (*ts*!*j*). $x \neq$ *vars-term-list* (*ts*!*j*) ! $i' \longrightarrow \sigma\ x =$ $\tau\ x$
    unfolding *ts*(*2*) **using** *ts*(*1*) **by** *auto*
  with *Cons*(*2*) i' *ts* **have** *IH*:*ctxt-of-pos-term p* ((*ts*!*j*) $\cdot \sigma$) = *ctxt-of-pos-term p* ((*ts*!*j*) $\cdot \tau$)
    using *Cons*(*1*)[*of ts*!*j* i'] **by** (*meson linear-term.simps*(*2*) *nth-mem*)
  {**fix** j' **assume** j':$j' <$ *length ts* $j' \neq j$
    with *Cons*(*2*) **have** *vars-term* (*ts* ! j') $\cap$ *vars-term* (*ts*!*j*) = {}
      unfolding *ts*(*2*) **by** (*metis disjoint-iff ts*(*1*) *var-in-linear-args*)
    then have $\forall x \in$ *vars-term* (*ts*!j'). $\sigma\ x = \tau\ x$
      using *Cons*(*5*) j' $*$ **by** (*metis disjoint-iff* i'(*1*) *length-var-poss-list nth-mem set-vars-term-list term.set-intros*(*4*) *ts*(*2*))
    then have (*ts*!j') $\cdot \sigma$ = (*ts*!j') $\cdot \tau$
      by (*meson term-subst-eq*)
  } **note** t'=*this*
  with *ts*(*1*) **have** *take j* (*map* ($\lambda t.\ t \cdot \sigma$) *ts*) = *take j* (*map* ($\lambda t.\ t \cdot \tau$) *ts*)
    using *nth-take-lemma*[*of j* (*map* ($\lambda t.\ t \cdot \sigma$) *ts*) (*map* ($\lambda t.\ t \cdot \tau$) *ts*)] **by** *simp*
  moreover from t' *ts*(*1*) **have** (*drop* (*Suc j*) (*map* ($\lambda t.\ t \cdot \sigma$) *ts*)) = (*drop* (*Suc*

j) (*map* ($\lambda t.\ t \cdot \tau$) *ts*))
using *nth-drop-equal*[*of* (*map* ($\lambda t.\ t \cdot \sigma$) *ts*) (*map* ($\lambda t.\ t \cdot \tau$) *ts*) *Suc j*] **by** *auto*
ultimately show *?case*
unfolding *ts*(*2*) *eval-term.simps ctxt-of-pos-term.simps* **using** *IH* **by** (*simp add*: *ts*(*1*))
qed *simp*

lemma *ctxt-apply-ctxt-apply*:
assumes $p \in$ *poss t*
shows (*ctxt-of-pos-term* (*p*@*q*) ((*ctxt-of-pos-term p t*) $\langle s \rangle$)) $\langle u \rangle$ = (*ctxt-of-pos-term p t*)$\langle$(*ctxt-of-pos-term q s*) $\langle u \rangle\rangle$
by (*metis assms ctxt-ctxt ctxt-of-pos-term-append hole-pos-ctxt-of-pos-term hole-pos-id-ctxt hole-pos-poss replace-at-subt-at*)

lemma *replace-at-append-subst*:
assumes *linear-term t*
and $p \in$ *poss t t*|-*p* = *Var x*
shows (*ctxt-of-pos-term* (*p*@*q*) ($t \cdot \sigma$)) $\langle s \rangle$ = $t \cdot \sigma$(*x* := (*ctxt-of-pos-term q* (σ *x*)) $\langle s \rangle$)
using *assms* **proof**(*induct p arbitrary*:*t*)
case (*Cons i p*)
then obtain *f ts* **where** *t*:*t* = *Fun f ts* **and** *i*:$i <$ *length ts* **and** *p*:$p \in$ *poss* (*ts*!*i*)
by (*meson args-poss*)
from *Cons*(*4*) **have** *x*:(*ts*!*i*)|-*p* = *Var x*
unfolding *t* **by** *simp*
from *Cons*(*2*) **have** *lin*:*linear-term* (*ts*!*i*)
using *i t* **by** *simp*
have *IH*:(*ctxt-of-pos-term* (*p*@*q*) ((*ts*!*i*) $\cdot\ \sigma$)) $\langle s \rangle$ = (*ts*!*i*) $\cdot\ \sigma$(*x* := (*ctxt-of-pos-term q* (σ *x*)) $\langle s \rangle$)
using *Cons*(*1*)[*OF lin p x*] **.**
let *?σ*=σ(*x* := (*ctxt-of-pos-term q* (σ *x*)) $\langle s \rangle$)
{fix *j* **assume** *j*:$j <$ *length ts* $j \neq i$
from *x* **have** $x \in$ *vars-term* (*ts*!*i*)
by (*metis p subsetD term.set-intros*(*3*) *vars-term-subt-at*)
then have $x \notin$ *vars-term* (*ts*!*j*)
using *j Cons*(*2*) **unfolding** *t* **by** (*meson i var-in-linear-args*)
then have (*ts*!*j*) $\cdot\ \sigma$ = (*ts*!*j*) $\cdot$ *?σ*
by (*simp add*: *term-subst-eq-conv*)
} **note** *sigma*=*this*
then have *take i* (*map* ($\lambda t.\ t \cdot \sigma$) *ts*) = *take i* (*map* ($\lambda t.\ t \cdot$ *?σ*) *ts*)
using *nth-take-lemma*[*of i* (*map* ($\lambda t.\ t \cdot \sigma$) *ts*) (*map* ($\lambda t.\ t \cdot$ *?σ*) *ts*)] *i* **by** *simp*
moreover from *sigma* **have** *drop* (*Suc i*) (*map* ($\lambda t.\ t \cdot \sigma$) *ts*) = *drop* (*Suc i*) (*map* ($\lambda t.\ t \cdot$ *?σ*) *ts*)
using *nth-drop-equal*[*of* (*map* ($\lambda t.\ t \cdot \sigma$) *ts*) (*map* ($\lambda t.\ t \cdot$ *?σ*) *ts*)] *i* **by** *simp*
ultimately show *?case*
unfolding *t append-Cons eval-term.simps ctxt-of-pos-term.simps intp-actxt.simps nth-map*[*OF i*] *IH*
by (*metis i id-take-nth-drop length-map nth-map*)
qed *simp*

lemma *replace-at-fun-poss-not-below*:
assumes $\neg\ p \leq_p q$
and $p \in$ *poss t* **and** $q \in$ *fun-poss* (*replace-at t p s*)
shows $q \in$ *fun-poss t*
using *assms* **by** (*metis ctxt-supt-id fun-poss-ctxt-apply-term hole-pos-ctxt-of-pos-term less-eq-pos-def*)

lemma *substitution-subterm-at*:
assumes $\forall j <$ *length* (*vars-term-list l*). σ (*vars-term-list l* ! *j*) = *s* |- (*var-poss-list l* ! *j*)
and $\exists \tau.\ l \cdot \tau = s$
shows $l \cdot \sigma = s$
using *assms* **proof**(*induct l arbitrary*:*s*)
case (*Var x*)
then show *?case*
unfolding *vars-term-list.simps poss-list.simps var-poss.simps eval-term.simps* **by** *simp*
next
case (*Fun f ts*)
from *Fun*(*3*) **obtain** *ss* **where** *s*:*s* = *Fun f ss* **and** *l*:*length ts* = *length ss*
by *fastforce*
{**fix** *i* **assume** *i*:*i* < *length ts*
{**fix** *j* **assume** *j*:*j*<*length* (*vars-term-list* (*ts*!*i*))
let *?p*=*var-poss-list* (*ts*!*i*) ! *j*
let *?x*=*vars-term-list* (*ts*!*i*) ! *j*
let *?k*=*sum-list* (*map* (*length* ∘ *vars-term-list*) (*take i ts*)) + *j*
from *i j* **have** *x*:*?x* = *vars-term-list* (*Fun f ts*) ! *?k*
unfolding *vars-term-list.simps* **by** (*simp add*: *concat-nth take-map*)
have *p*:*var-poss-list* (*Fun f ts*) ! *?k* = *i* # *?p* **proof**−
from *i* **have** *i'*:*i* < *length* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length ts*] (*map var-poss-list ts*))
by *simp*
from *i j* **have** *j* < *length* ((*map var-poss-list ts*) ! *i*)
using *length-var-poss-list* **by** (*metis* (*mono-tags*, *lifting*) *nth-map*)
with *i* **have** *j'*:*j* < *length* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length ts*] (*map var-poss-list ts*) ! *i*)
by *simp*
{**fix** *l* **assume** *l* < *length ts*
then have (*map length* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length ts*] (*map var-poss-list ts*)))!*l* = (*map* (*length* ∘ *vars-term-list*) *ts*) ! *l*
using *length-var-poss-list* **by** *simp*
}
then have *map length* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length ts*] (*map var-poss-list ts*)) = *map* (*length* ∘ *vars-term-list*) *ts*
using *nth-equalityI*[**where** *ys*=*map* (*length* ∘ *vars-term-list*) *ts*] **by** *simp*
with *i* **have** *k*:*sum-list* (*map length* (*take i* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length ts*] (*map var-poss-list ts*)))) + *j* = *?k*
by (*metis take-map*)

then have *var-poss-list* (*Fun f ts*) ! *?k* = (*map2* (λ*i*. *map* ((#) *i*)) [*0*..<*length ts*] (*map var-poss-list ts*))!*i* !*j*
unfolding *var-poss-list.simps* **using** *concat-nth*[*OF i' j'*] **by** *presburger*
also have ... = (*map* ((#) *i*) (*var-poss-list* (*ts*!*i*)))!*j* **using** *i* **by** *simp*
also have ... = *i* # *?p* **using** *nth-map j length-var-poss-list* **by** *metis*
ultimately show *?thesis* **by** *simp*
qed
from *i j* **have** *k*:*?k* < *length* (*vars-term-list* (*Fun f ts*))
unfolding *vars-term-list.simps* **by** (*metis concat-nth-length length-map map-map nth-map take-map*)
from *Fun*(*2*) *k* **have** σ *?x* = (*ss*!*i*) |- (*var-poss-list* (*ts*!*i*) ! *j*)
unfolding *x s* **using** *p* **by** *simp*
}
then have ∀ *j*<*length* (*vars-term-list* (*ts*!*i*)).σ (*vars-term-list* (*ts*!*i*) ! *j*) = (*ss*!*i*) |- *var-poss-list* (*ts*!*i*) ! *j*
by *simp*
moreover from *Fun*(*3*) **have** ∃ τ. (*ts*!*i*) · τ = *ss*!*i*
unfolding *eval-term.simps s* **using** *i l* **by** (*metis nth-map term.inject*(*2*))
ultimately have (*ts*!*i*) · σ = *ss*!*i*
using *i Fun*(*1*) *nth-mem* **by** *blast*
}
then show *?case* **unfolding** *eval-term.simps s*
using *l* **by** (*simp add*: *map-nth-eq-conv*)
qed

lemma *vars-map-vars-term*:
map f (*vars-term-list t*) = *vars-term-list* (*map-vars-term f t*)
unfolding *map-vars-term-eq* **proof**(*induct t*)
case (*Fun g ts*)
then have *map* (λ*xs*. *map f xs*)(*map vars-term-list ts*) = *map vars-term-list* (*map* (λ*t*. *t* · (*Var* ∘ *f*)) *ts*)
by *fastforce*
then show *?case* **unfolding** *vars-term-list.simps eval-term.simps map-map map-concat*
by *presburger*
qed (*simp add*: *vars-term-list.simps*)

lemma *ctxt-apply-subt-at*:
assumes *q* ∈ *poss s*
shows (*ctxt-of-pos-term p* (*s*|-*q*)) ⟨*t*⟩ = (*ctxt-of-pos-term* (*q*@*p*) *s*) ⟨*t*⟩ |- *q*
using *assms* **proof**(*induct q arbitrary*: *s*)
case (*Cons i q*)
from *Cons*(*2*) **obtain** *f ss* **where** *i*:*i* < *length ss* **and** *s*:*s* = *Fun f ss*
by (*meson args-poss*)
from *i Cons* **show** *?case* **unfolding** *s*
by (*metis ctxt-apply-ctxt-apply ctxt-supt-id replace-at-subt-at*)
qed *simp*

1.3.1 Utilities for *mk-subst*

We consider the special case of applying *mk-subst* when the variables involved form a partition.

lemma *mk-subst-same*:
  assumes *length xs* = *length ts distinct xs*
  shows *map* (*mk-subst f* (*zip xs ts*)) *xs* = *ts*
  using *assms* **by** (*simp add*: *mk-subst-distinct map-nth-eq-conv*)

lemma *map2-zip*: *set* (*map fst* (*concat* (*map2 zip xs ys*))) ⊆ *set* (*concat xs*)
proof
  fix *x* **assume** *x*:*x* ∈ *set* (*map fst* (*concat* (*map2 zip xs ys*)))
  let *?l*=*min* (*length xs*) (*length ys*)
  from *x* **obtain** *i* **where** *i*:*i* < *?l x* ∈ *set* (*map fst* (*zip* (*xs*!*i*) (*ys*!*i*)))
   by (*smt* (*verit*) *case-prod-conv in-set-conv-nth length-map length-zip min.strict-boundedE nth-concat-split nth-map nth-zip*)
  then have *x* ∈ *set* (*xs*!*i*)
    by (*metis in-set-takeD map-fst-zip-take*)
  then show *x* ∈ *set* (*concat xs*)
   using *i*(*1*) **by** (*metis concat-nth concat-nth-length in-set-conv-nth min.strict-boundedE*)

qed

lemma *mk-subst-partition*:
  fixes *xs* :: *'a list list*
  assumes *l*:*length xs* = *length ss*
    and *part*:*is-partition* (*map set xs*)
   shows ∀ *i* < *length xs*. ∀ *x* ∈ *set* (*xs*!*i*). (*mk-subst f* (*zip* (*xs*!*i*) (*ss*!*i*))) *x* = (*mk-subst f* (*concat* (*map2 zip xs ss*))) *x*
proof−
  {**fix** *i* **assume** *i*:*i* < *length xs*
    {**fix** *x* **assume** *x*:*x* ∈ *set* (*xs*!*i*)
      have *concat* (*map2 zip xs ss*) = *concat* (*map2 zip* (*take i xs*) (*take i ss*)) @ *concat* (*map2 zip* (*drop i xs*) (*drop i ss*))
        by (*metis append-take-drop-id concat-append drop-map drop-zip take-map take-zip*)
      moreover have *concat* (*map2 zip* (*drop i xs*) (*drop i ss*)) = *concat* (*zip* (*xs*!*i*) (*ss*!*i*) # (*map2 zip* (*drop* (*Suc i*) *xs*) (*drop* (*Suc i*) *ss*)))
      using *i l* **by** (*smt* (*verit*, *del-insts*) *Cons-nth-drop-Suc list.map*(*2*) *prod.simps*(*2*) *zip-Cons-Cons*)
      ultimately have *cc*:*concat* (*map2 zip xs ss*) = *concat* (*map2 zip* (*take i xs*) (*take i ss*)) @
              concat (*zip* (*xs*!*i*) (*ss*!*i*) # (*map2 zip* (*drop* (*Suc i*) *xs*) (*drop* (*Suc i*) *ss*))) **by** *presburger*
      {**fix** *j* **assume** *j* < *length xs* **and** *j* ≠ *i*
        with *i x part* **have** *x* ∉ *set* (*xs*!*j*)
          unfolding *is-partition-alt is-partition-alt-def* **by** *auto*
      } **note** *part*=*this*
      then have *x* ∉ *set* (*concat* (*take i xs*))

by (*smt* (*verit*) *in-set-conv-nth length-take less-length-concat min.strict-boundedE nth-take order-less-irrefl*)
then have $x \notin$ *set* (*map fst* (*concat* (*map2 zip* (*take i xs*) (*take i ss*))))
using *map2-zip in-mono* **by** *fastforce*
then have *subst*:(*mk-subst f* (*concat* (*map2 zip xs ss*))) x = (*mk-subst f* (*concat* (*zip* (*xs*!*i*) (*ss*!*i*) #
(*map2 zip* (*drop* (*Suc i*) *xs*) (*drop* (*Suc i*) *ss*))))) x
unfolding *cc* **using** *mk-subst-concat* **by** *metis*
have *mk-subst f* (*zip* (*xs* ! *i*) (*ss* ! *i*)) x = (*mk-subst f* (*concat* (*map2 zip xs ss*))) x
proof(*cases* $x \in$ *set* (*map fst* (*zip* (*xs*!*i*) (*ss*!*i*))))
case *True*
then show *?thesis*
using *mk-subst-concat-Cons subst* **by** *metis*
next
case *False*
{**fix** *j* **assume** *j*:$j <$ *length* (*drop* (*Suc i*) *xs*)
then have (*drop* (*Suc i*) *xs*)!*j* = *xs*!(*Suc i* + *j*)
using *Suc-leI i nth-drop* **by** *blast*
moreover from *i j* **have** *Suc i* + *j* < *length xs*
by (*metis add.commute length-drop less-diff-conv*)
ultimately have $x \notin$ *set* ((*drop* (*Suc i*) *xs*)!*j*)
using *part* **by** (*metis Suc-n-not-le-n le-add1*)
}
then have $x \notin$ *set* (*concat* (*drop* (*Suc i*) *xs*))
by (*smt* (*verit*) *in-set-conv-nth length-map length-take less-not-refl2 min.strict-boundedE nth-concat-split nth-map nth-take*)
then have $x \notin$ *set* (*map fst* (*concat* (*map2 zip* (*drop* (*Suc i*) *xs*) (*drop* (*Suc i*) *ss*))))
using *map2-zip in-mono* **by** *fastforce*
with *False* **have** $x \notin$ *set* (*map fst* (*concat* (*zip* (*xs*!*i*) (*ss*!*i*) # (*map2 zip* (*drop* (*Suc i*) *xs*) (*drop* (*Suc i*) *ss*)))))
unfolding *concat.simps* **by** (*metis Un-iff map-append set-append*)
with *False* **show** *?thesis*
unfolding *subst* **using** *mk-subst-not-mem'* **by** *metis*
qed
}
}
then show *?thesis* **by** *simp*
qed

The following lemma is used later to show that $A = (\textit{to-pterm}\ (\textit{lhs}\ \alpha)) \cdot \sigma$ implies $A = (\textit{to-pterm}\ (\textit{lhs}\ \alpha)) \cdot \langle As\rangle_\alpha$ for some suitable As.

lemma *subst-imp-mk-subst*:
assumes $s = t \cdot \sigma$
shows $\exists$ *ss*. $t \cdot \sigma = t \cdot$ (*mk-subst Var* (*zip* (*vars-distinct t*) *ss*)) $\wedge$ *length ss* = *length* (*vars-distinct t*) $\wedge$ ($\forall$ $i <$ *length ss*. σ (*vars-distinct t*!*i*) = *ss*!*i*)
proof−

let *?ss=map σ (vars-distinct t)*
let *?τ=(mk-subst Var (zip (vars-distinct t) ?ss))*
{fix *x* **assume** *x ∈ vars-term t*
then have *σ x = ?τ x* **unfolding** *mk-subst-def*
by *(simp add: map-of-zip-map)*
}
then have *t · σ = t · ?τ*
using *term-subst-eq* **by** *blast*
then show *?thesis* **by** *auto*
qed

lemma *mk-subst-rename*:
assumes *length (vars-distinct t) = length xs* **and** *inj f*
shows *t · (mk-subst Var (zip (vars-distinct t) xs)) = (map-vars-term f t) · (mk-subst Var (zip (vars-distinct (map-vars-term f t)) xs))*
proof−
{fix *x* **assume** *x ∈ vars-term t*
then obtain *i* **where** *i:x = (vars-distinct t)!i i < length (vars-distinct t)*
by *(metis in-set-conv-nth set-vars-term-list vars-term-list-vars-distinct)*
with *assms* **have** *1:(mk-subst Var (zip (vars-distinct t) xs)) x = xs!i*
using *mk-subst-distinct* **by** *(metis comp-apply distinct-remdups distinct-rev)*
have *vars-distinct (map-vars-term f t) = map f (vars-distinct t)*
unfolding *vars-map-vars-term[symmetric] comp-apply* **using** *assms(2)*
by *(metis distinct-map distinct-remdups distinct-remdups-id inj-on-inverseI remdups-map-remdups rev-map the-inv-f-f)*
with *assms i* **have** *2:(mk-subst Var (zip (vars-distinct (map-vars-term f t)) xs)) (f x) = xs!i*
by *(metis (mono-tags, lifting) comp-apply distinct-remdups distinct-rev length-map mk-subst-same nth-map)*
from *1 2* **have** *(mk-subst Var (zip (vars-distinct t) xs)) x = (mk-subst Var (zip (vars-distinct (map-vars-term f t)) xs)) (f x)*
by *presburger*
}
then show *?thesis*
by *(simp add: apply-subst-map-vars-term term-subst-eq-conv)*
qed

1.4 Matching Terms

The goal is showing that *match* $(t \cdot \sigma)$ t = *Some* σ whenever the domain of σ is a subset of the variables in t. For that we need some helper lemmas.

lemma *decompose-fst*:
assumes *decompose (Fun f ss) t = Some us*
shows *map fst us = ss*
proof−
from *assms* **obtain** *ts* **where** *t:t = Fun f ts*
by *(metis (no-types, lifting) decompose-def option.distinct(1) decompose-Some is-FunE old.prod.case term.case-eq-if)*
with *assms* **have** *length ss = length ts*

by *blast*
with *assms*(*1*) *t* **show** *?thesis*
by *auto*
qed

lemma *decompose-vars-term*:
assumes *decompose* (*Fun f ss*) *t* = *Some us*
shows *vars-term* (*Fun f ss*) = ($\bigcup$(*a*, *b*) ∈ *set us*. *vars-term a*)
proof−
have *vars-term* (*Fun f ss*) = ($\bigcup$ *s* ∈ *set ss*. *vars-term s*)
by (*meson Term.term.simps*(*18*))
also have ... = ($\bigcup$ *s* ∈ *set* (*map fst us*). *vars-term s*)
using *assms decompose-fst* **by** *metis*
finally show *?thesis*
using *image-image* **by** *auto*
qed

lemma *match-term-list-domain*:
assumes *match-term-list P σ* = *Some τ*
shows ∀ *x*. *x* ∉ ($\bigcup$ (*a*, *b*) ∈ *set P*. *vars-term a*) ∧ *σ x* = *None* ⟶ *τ x* = *None*
using *assms* **proof**(*induct P σ rule*:*match-term-list.induct*)
case (*2 x t P σ*)
then show *?case*
by (*metis* (*mono-tags*, *lifting*) *Sup-insert Un-iff case-prod-conv fun-upd-idem-iff fun-upd-triv fun-upd-twist image-insert list.simps*(*15*) *match-term-list.simps*(*2*) *option.simps*(*3*) *term.set-intros*(*3*))
next
case (*3 f ss g ts P σ*)
from *3*(*2*) **obtain** *us* **where** *us*:*decompose* (*Fun f ss*) (*Fun g ts*) = *Some us*
using *match-term-list.simps*(*3*) *option.distinct*(*1*) *option.simps*(*4*) **by** *fastforce*
with *3*(*2*) **have** ∗:*match-term-list* (*us* @ *P*) *σ* = *Some τ*
by *auto*
from *us* **have** ($\bigcup$ (*a*, *b*) ∈ *set* ((*Fun f ss*, *Fun g ts*) # *P*). *vars-term a*) = ($\bigcup$ *s* ∈ *set ss*. *vars-term s*) ∪ ($\bigcup$ (*a*, *b*) ∈ *set P*. *vars-term a*)
by *simp*
also have ... = ($\bigcup$ (*a*, *b*) ∈ *set* (*us*@*P*). *vars-term a*)
using *us* **by** (*metis* (*mono-tags*, *lifting*) *Term.term.simps*(*18*) *UN-Un decompose-vars-term set-append*)
finally show *?case*
using *3*(*1*) *us* ∗ **by** *metis*
qed *simp-all*

lemma *match-subst-domain*:
assumes *match a b* = *Some σ*
shows *subst-domain σ* ⊆ *vars-term b*
proof−
from *assms* **have** ∀ *x*. *x* ∉ *vars-term b* ⟶ *σ x* = *Var x*
unfolding *match-def match-list-def subst-of-map-def* **using** *match-term-list-domain* **by** *fastforce*

then show *?thesis*
using *subst-domain-def* **by** *fastforce*
qed

lemma *match-trivial*:
assumes $subst\text{-}domain\ \sigma \subseteq vars\text{-}term\ t$
shows $match\ (t \cdot \sigma)\ t = Some\ \sigma$
proof−
obtain τ **where** $tau{:}match\ (t \cdot \sigma)\ t = Some\ \tau$ **and** $1{:}(\forall x{\in}vars\text{-}term\ t.\ \sigma\ x = \tau\ x)$
by (*meson match-complete′*)
from *assms* **have** $2{:}\forall x.\ x \notin vars\text{-}term\ t \longrightarrow \sigma\ x = Var\ x$
by (*meson notin-subst-domain-imp-Var subset-eq*)
from *tau* **have** $3{:}\forall x.\ x \notin vars\text{-}term\ t \longrightarrow \tau\ x = Var\ x$
using *match-subst-domain notin-subst-domain-imp-Var* **by** *fastforce*
from *1 2 3* **show** *?thesis*
by (*metis subst-term-eqI tau term-subst-eq*)
qed

end

1.4.1 Matching of Linear Terms

theory *Linear-Matching*
imports *Proof-Term-Utils*
begin

For a linear term the matching substitution can simply be computed with the following definition.

definition $match\text{-}substs :: ('f, 'v)\ term \Rightarrow ('f, 'v)\ term \Rightarrow ('v \times ('f, 'v)\ term)\ list$
where $match\text{-}substs\ t\ s = (zip\ (vars\text{-}term\text{-}list\ t)\ (map\ (\lambda p.\ s|\text{-}p)\ (var\text{-}poss\text{-}list\ t)))$

lemma *mk-subst-partition-special*:
assumes $length\ ss = length\ ts$
and $is\text{-}partition\ (map\ vars\text{-}term\ ts)$
shows $\forall i < length\ ts.\ (ts!i) \cdot (mk\text{-}subst\ f\ (zip\ (vars\text{-}term\text{-}list\ (ts!i))\ (ss!i))) = (ts!i) \cdot (mk\text{-}subst\ f\ (concat\ (map2\ zip\ (map\ vars\text{-}term\text{-}list\ ts)\ ss)))$
proof−
let $?xs{=}map\ vars\text{-}term\text{-}list\ ts$
have $xs{:}map\ vars\text{-}term\ ts = map\ set\ (map\ vars\text{-}term\text{-}list\ ts)$ **by** *simp*
from *assms*(*1*) **have** $l{:}length\ ?xs = length\ ss$ **by** *simp*
{**fix** i **assume** $i{:}i < length\ ts$
{**fix** x **assume** $x \in vars\text{-}term\ (ts!i)$
then have $mk\text{-}subst\ f\ (zip\ (map\ vars\text{-}term\text{-}list\ ts\ !\ i)\ (ss\ !\ i))\ x = mk\text{-}subst\ f\ (concat\ (map2\ zip\ (map\ vars\text{-}term\text{-}list\ ts)\ ss))\ x$
using *i mk-subst-partition*[*OF l assms*(*2*)[*unfolded xs*]] **by** *simp*
}
then have $ts!i \cdot (mk\text{-}subst\ f\ (zip\ (vars\text{-}term\text{-}list\ (ts!i))\ (ss!i))) = (ts!i) \cdot$

(*mk-subst f* (*concat* (*map2 zip* (*map vars-term-list ts*) *ss*)))
    by (*simp add*: *i term-subst-eq-conv*)
  }
  then show *?thesis* **by** *fastforce*
qed

lemma *match-substs-Fun*:
  assumes *l*:*length ts* = *length ss*
  shows *match-substs* (*Fun f ts*) (*Fun g ss*) = *concat* (*map2 zip* (*map vars-term-list ts*) (*map2* (λ*t s*. *map* ((|-) *s*) (*var-poss-list t*)) *ts ss*))
    (**is** *match-substs* (*Fun f ts*) (*Fun g ss*) = *concat* (*map2 zip ?xs ?terms*))
proof−
  have *l′*:*length ?xs* = *length ?terms* **using** *l* **by** *simp*
  {fix *i* **assume** *i* < *length ?xs*
    then have *i*:*i* < *length ts* **by** *simp*
    with *l* **have** *zip*:(*zip ts ss*)!*i* = (*ts*!*i*, *ss*!*i*) **by** *simp*
    from *i l* **have** *length* (*map vars-term-list ts* ! *i*) = *length* (*map* (λ*p*. (*ss*!*i*)|-*p*) (*var-poss-list* (*ts*!*i*)))
      by (*simp add*: *length-var-poss-list*)
    with *zip* **have** *length* (*?xs*!*i*) = *length* (*?terms*!*i*)
      using *i l′* **by** *auto*
  }note *l-i*=*this*
  have *vars-term-list* (*Fun f ts*) = *concat ?xs*
    unfolding *vars-term-list.simps* **by** *simp*
  moreover have *map* ((|-) (*Fun g ss*)) (*var-poss-list* (*Fun f ts*)) = *concat ?terms*
proof−
    have *l-map2*:*length* (*map2* (λ*i*. *map* ((#) *i*)) [*0*..<*length ts*] (*map var-poss-list ts*)) = *length ts*
      unfolding *length-map length-zip* **by** *simp*
    {fix *i* **assume** *i*:*i* < *length ts*
     with *l* **have** *length* (*map2* (λ*i*. *map* ((#) *i*)) [*0*..<*length ts*] (*map var-poss-list ts*) !*i*) = *length* (*map var-poss-list ts*!*i*)
        unfolding *nth-map*[*OF i*] **by** *simp*
    }
     with *l-map2* **have** *length* (*map* ((|-) (*Fun g ss*)) (*var-poss-list* (*Fun f ts*))) = *length* (*concat* (*map var-poss-list ts*))
    unfolding *var-poss-list.simps length-map length-concat* **by** (*smt* (*verit*, *del-insts*) *length-map map-nth-eq-conv*)
   moreover have *length* (*concat ?terms*) = *length* (*concat* (*map var-poss-list ts*))
proof−
      {fix *i* **assume** *i* < *length ts*
        with *l* **have** *length* (*map2* (λ*t s*. *map* ((|-) *s*) (*var-poss-list t*)) *ts ss* ! *i*) = *length* (*map var-poss-list ts*!*i*) **by** *simp*
      }
       moreover have *length* (*map2* (λ*t s*. *map* ((|-) *s*) (*var-poss-list t*)) *ts ss*) = *length ts* **using** *l* **by** *simp*
     ultimately show *?thesis* **unfolding** *length-concat* **by** (*smt* (*verit*, *del-insts*) *length-map map-nth-eq-conv*)
    qed

ultimately have *l''*:*length* (*map* ((|-) (*Fun g ss*)) (*var-poss-list* (*Fun f ts*))) = *length* (*concat ?terms*) **by** *presburger*
{fix *i* **assume** *i*:*i* < *length* (*var-poss-list* (*Fun f ts*))
let *?ps*=*map2* (λ*i*. *map* ((#) *i*)) [*0*..<*length ts*] (*map var-poss-list ts*)
let *?p*=*var-poss-list* (*Fun f ts*) ! *i*
from *l* **have** *l-terms*:*length ?terms* = *length ts* **by** *auto*
from *l* **have** *l-ps*:*length ?ps* = *length ts* **by** *auto*
obtain *j k* **where** *j*:*j* < *length ts* **and** *k*:*k* < *length* (*var-poss-list* (*ts*!*j*)) **and** *i-sum-list*:*i* = *sum-list* (*map length* (*take j ?ps*)) + *k*
and ∗:*?p* = *map2* (λ*i*. *map* ((#) *i*)) [*0*..<*length ts*] (*map var-poss-list ts*) ! *j* ! *k*
using *less-length-concat*[*OF i*[*unfolded var-poss-list.simps*]] **by** *auto*
let *?p'*=(*var-poss-list* (*ts*!*j*)) ! *k*
from ∗ *k j* **have** *p*:*?p* = *j* # *?p'* **by** *simp*
from *j l* **have** *1*:(*Fun g ss*) |- *?p* = (*ss*!*j*) |- *?p'* **unfolding** *p* **by** *simp*
have *i-sum-list2*:*i* = *sum-list* (*map length* (*take j ?terms*)) + *k* **proof**−
{fix *n* **assume** *n* < *length ts*
with *l* **have** *length* (*?terms*!*n*) = *length* (*?ps*!*n*) **by** *auto*
}
then have *map length ?terms* = *map length ?ps*
using *l-terms l-ps* **by** (*simp add*: *map-eq-conv'*)
then show *?thesis* **unfolding** *i-sum-list* **by** (*metis take-map*)
qed
from *j k* **have** *k* < *length* (*?terms* ! *j*) **by** (*smt* (*verit*) *l-i length-map length-var-poss-list nth-map*)
with *j i-sum-list2* **have** *concat ?terms* ! *i* = *?terms* ! *j* ! *k*
using *concat-nth*[*of j ?terms k i*] **unfolding** *length-map length-zip l min.idem* **by** *auto*
then have *2*:*concat ?terms* ! *i* = (*ss*!*j*) |- *?p'* **using** *k j l* **by** *auto*
from *1 2* **have** *map* ((|-) (*Fun g ss*)) (*var-poss-list* (*Fun f ts*)) ! *i* = *concat ?terms* ! *i*
unfolding *var-poss-list.simps nth-map*[*OF i*[*unfolded var-poss-list.simps*]] **by** *simp*
}
with *l''* **show** *?thesis* **by** (*metis length-map nth-equalityI*)
qed
ultimately show *?thesis*
unfolding *match-substs-def* **using** *concat-map2-zip*[*OF l'*] *l-i* **by** *presburger*
qed

If all function symbols in term t coincide with function symbols in term s, then t matches s.

lemma *fun-poss-eq-imp-matches*:
fixes *s t* :: (*'f*, *'v*) *term*
assumes *linear-term t* **and** ∀ *p* ∈ *poss t*. ∀ *f ts*. *t*|-*p* = *Fun f ts* ⟶ (∃ *ss*. *length ts* = *length ss* ∧ *s*|-*p* = *Fun f ss*)
shows *t* · (*mk-subst Var* (*match-substs t s*)) = *s*
using *assms* **proof**(*induct t arbitrary*:*s*)
case (*Var x*)

have match-substs (Var x) s = [(x, s)]
unfolding match-substs-def var-poss-list.simps vars-term-list.simps **by** simp
then show ?case **by** simp
next
case (Fun f ts)
from Fun(3) **obtain** ss **where** l:length ts = length ss **and** s:s = Fun f ss **by** auto
let ?σ=mk-subst Var (match-substs (Fun f ts) (Fun f ss))
let ?xs=map vars-term-list ts
let ?ss=map (λ(t, s). map (λp. s|-p) (var-poss-list t)) (zip ts ss)
have concat-zip:match-substs (Fun f ts) (Fun f ss) = concat (map2 zip ?xs ?ss)
unfolding match-substs-Fun[OF l] **by** simp
from Fun(2) **have** part:is-partition (map set ?xs)
by (smt (verit, ccfv-SIG) length-map linear-term.elims(2) map-nth-eq-conv set-vars-term-list term.distinct(1) term.sel(4))
have l′:length ?xs = length ?ss **using** l **by** simp
{fix i **assume** i:i < length ts
with Fun(2) **have** lin:linear-term (ts!i) **by** simp
let ?σi=mk-subst Var (match-substs (ts!i) (ss!i))
{fix p f′ ts′ **assume** p:p ∈ poss (ts!i) ts!i |- p = Fun f′ ts′
from p(1) i **have** i#p ∈ poss (Fun f ts) **by** simp
moreover from p(2) i **have** (Fun f ts)|-(i#p) = Fun f′ ts′ **by** simp
ultimately obtain ss′ **where** length ts′ = length ss′ **and** s|-(i#p) = Fun f′ ss′ **using** Fun(3) **by** blast
then have ∃ ss′. length ts′ = length ss′ ∧ (ss!i)|-p = Fun f′ ss′ **unfolding** s **by** simp
}
then have ∀ p ∈ poss (ts!i). ∀ f′ ts′. (ts!i)|-p = Fun f′ ts′ ⟶ (∃ ss′. length ts′ = length ss′ ∧ (ss!i)|-p = Fun f′ ss′) **by** simp
with Fun(1) lin i **have** IH:(ts!i) · ?σi = ss!i **using** nth-mem **by** blast
have (ts!i) · ?σ = (ts!i) · ?σi **proof**−
{fix x **assume** x:x ∈ vars-term (ts!i)
from i l **have** ∗:map2 (λt s. map ((|-) s) (var-poss-list t)) ts ss ! i = map ((|-) (ss ! i)) (var-poss-list (ts ! i)) **by** auto
with i x **have** ?σ x = ?σi x
unfolding concat-zip **using** mk-subst-partition[OF l′ part] **unfolding** s match-substs-Fun[OF l] match-substs-def length-map
by (smt (verit, best) nth-map set-vars-term-list)
}
then show ?thesis **by** (simp add: term-subst-eq-conv)
qed
then have (ts!i) · ?σ = ss!i **using** IH i **by** presburger
}
then show ?case **unfolding** s **by** (simp add: l map-nth-eq-conv)
qed

end

2 Proof Terms

theory *Proof-Terms*
  imports
    First-Order-Terms.Matching
    First-Order-Rewriting.Multistep
    Proof-Term-Utils
begin

2.1 Definitions

A rewrite rule consists of a pair of terms representing its left-hand side and right-hand side. We associate a rule symbol with each rewrite rule.

datatype $('f, 'v)$ *prule* =
  Rule (*lhs*: $('f, 'v)$ *term*) (*rhs*: $('f, 'v)$ *term*) (- → - [*51*, *51*] *52*)

Translate between *prule* defined here and *rule* as defined in IsaFoR.

abbreviation *to-rule* :: $('f, 'v)$ *prule* ⇒ $('f, 'v)$ *rule*
  where *to-rule r* ≡ (*lhs r*, *rhs r*)

Proof terms are terms built from variables, function symbols, and rules.

type-synonym
  $('f, 'v)$ *pterm* = $(('f, 'v)$ *prule* + $'f, 'v)$ *term*
type-synonym
  $('f, 'v)$ *pterm-ctxt* = $(('f, 'v)$ *prule* + $'f, 'v)$ *ctxt*

We provides an easier notation for proof terms (avoiding *Inl* and *Inr*).

abbreviation *Prule* :: $('f, 'v)$ *prule* ⇒ $('f, 'v)$ *pterm list* ⇒ $('f, 'v)$ *pterm*
  where *Prule* α *As* ≡ *Fun* (*Inl* α) *As*
abbreviation *Pfun* :: $'f$ ⇒ $('f, 'v)$ *pterm list* ⇒ $('f, 'v)$ *pterm*
  where *Pfun f As* ≡ *Fun* (*Inr f*) *As*

Also for contexts.

abbreviation *Crule* :: $('f, 'v)$ *prule* ⇒ $('f, 'v)$ *pterm list* ⇒ $('f, 'v)$ *pterm-ctxt* ⇒ $('f, 'v)$ *pterm list* ⇒$('f, 'v)$ *pterm-ctxt*
  where *Crule* α *As1 C As2* ≡ *More* (*Inl* α) *As1 C As2*
abbreviation *Cfun* :: $'f$ ⇒ $('f, 'v)$ *pterm list* ⇒ $('f, 'v)$ *pterm-ctxt* ⇒ $('f, 'v)$ *pterm list* ⇒$('f, 'v)$ *pterm-ctxt*
  where *Cfun f As1 C As2* ≡ *More* (*Inr f*) *As1 C As2*

Case analysis on proof terms.

lemma *pterm-cases* [*case-names Var Pfun Prule*, *cases type*: *pterm*]:
  $(\bigwedge x.\ A = Var\ x \Longrightarrow P) \Longrightarrow (\bigwedge f\ As.\ A = Pfun\ f\ As \Longrightarrow P) \Longrightarrow (\bigwedge \alpha\ As.\ A = Prule\ \alpha\ As \Longrightarrow P) \Longrightarrow P$
proof (*cases A*)
  case (*Fun x21 x22*)
  show $\bigwedge x21\ x22.\ (\bigwedge x.\ A = Var\ x \Longrightarrow P) \Longrightarrow (\bigwedge f\ As.\ A = Pfun\ f\ As \Longrightarrow P) \Longrightarrow (\bigwedge \alpha\ As.\ A = Prule\ \alpha\ As \Longrightarrow P) \Longrightarrow A = Fun\ x21\ x22 \Longrightarrow P$

using *sum.exhaust* **by** *auto*
qed

Induction scheme for proof terms.

lemma
fixes P :: $('f, 'v)$ *pterm* $\Rightarrow$ *bool*
assumes $\bigwedge x.\ P\ (Var\ x)$
and $\bigwedge f\ As.\ (\bigwedge a.\ a \in set\ As \Longrightarrow P\ a) \Longrightarrow P\ (Pfun\ f\ As)$
and $\bigwedge \alpha\ As.\ (\bigwedge a.\ a \in set\ As \Longrightarrow P\ a) \Longrightarrow P\ (Prule\ \alpha\ As)$
shows *pterm-induct* [*case-names Var Pfun Prule*, *induct type*: *pterm*]: $P\ A$
using *assms* **proof**(*induct A*)
case (*Fun f ts*)
then show *?case* **by**(*cases f*) *auto*
qed *simp*

Induction scheme for contexts of proof terms.

lemma
fixes P :: $('f, 'v)$ *pterm-ctxt* $\Rightarrow$ *bool*
assumes $P\ \Box$
and $\bigwedge f\ ss1\ C\ ss2.\ P\ C \Longrightarrow P\ (Cfun\ f\ ss1\ C\ ss2)$
and $\bigwedge \alpha\ ss1\ C\ ss2.\ P\ C \Longrightarrow P\ (Crule\ \alpha\ ss1\ C\ ss2)$
shows *pterm-ctxt-induct* [*case-names Hole Cfun Crule*, *induct type*: *pterm-ctxt*]:
$P\ C$
using *assms* **proof**(*induct C*)
case (*More f ss1 C ss2*)
then show *?case* **by**(*cases f*) *auto*
qed

Obtain the distinct variables occurring on the left-hand side of a rule in the order they appear.

abbreviation *var-rule* :: $('f, 'v)$ *prule* $\Rightarrow$ $'v$ *list* **where** *var-rule* $\alpha \equiv$ *vars-distinct* (*lhs* α)

abbreviation *lhs-subst* :: $('g, 'v)$ *term list* $\Rightarrow$ $('f, 'v)$ *prule* $\Rightarrow$ $('g, 'v)$ *subst* ($\langle$-$\rangle$-[*0*,*99*])
where *lhs-subst As* $\alpha \equiv$ *mk-subst Var* (*zip* (*var-rule* α) *As*)

A proof term using only function symbols and variables is an empty step.

fun *is-empty-step* :: $('f, 'v)$ *pterm* $\Rightarrow$ *bool* **where**
is-empty-step (*Var x*) = *True*
| *is-empty-step* (*Pfun f As*) = *list-all is-empty-step As*
| *is-empty-step* (*Prule* α *As*) = *False*

fun *is-Prule* :: $('f, 'v)$ *pterm* $\Rightarrow$ *bool* **where**
is-Prule (*Prule* - -) = *True*
| *is-Prule* - = *False*

Source and target

fun *source* :: $('f, 'v)$ *pterm* $\Rightarrow$ $('f, 'v)$ *term* **where**
source $(Var\ x) = Var\ x$
$|$ *source* $(Pfun\ f\ As) = Fun\ f\ (map\ source\ As)$
$|$ *source* $(Prule\ \alpha\ As) = lhs\ \alpha \cdot \langle map\ source\ As\rangle_\alpha$

fun *target* :: $('f, 'v)$ *pterm* $\Rightarrow$ $('f, 'v)$ *term* **where**
target $(Var\ x) = Var\ x$
$|$ *target* $(Pfun\ f\ As) = Fun\ f\ (map\ target\ As)$
$|$ *target* $(Prule\ \alpha\ As) = rhs\ \alpha \cdot \langle map\ target\ As\rangle_\alpha$

Source also works for proof term contexts in left-linear TRSs.

fun *source-ctxt* :: $('f, 'v)$ *pterm-ctxt* $\Rightarrow$ $('f, 'v)$ *ctxt* **where**
source-ctxt $\Box = \Box$
$|$ *source-ctxt* $(Cfun\ f\ As1\ C\ As2) = More\ f\ (map\ source\ As1)\ (source\text{-}ctxt\ C)\ (map\ source\ As2)$
$|$ *source-ctxt* $(Crule\ \alpha\ As1\ C\ As2) =$
$(let\ ctxt\text{-}pos = (var\text{-}poss\text{-}list\ (lhs\ \alpha))!(length\ As1);$
$placeholder = Var\ ((vars\text{-}term\text{-}list\ (lhs\ \alpha))\ !\ (length\ As1))\ in$
$ctxt\text{-}of\text{-}pos\text{-}term\ (ctxt\text{-}pos)\ (lhs\ \alpha \cdot \langle map\ source\ (As1\ @\ ((placeholder\ \#\ As2)))\rangle_\alpha))$
$\circ_c\ (source\text{-}ctxt\ C)$

abbreviation *co-initial* $A\ B \equiv (source\ A = source\ B)$

Transform simple terms to proof terms.

fun *to-pterm* :: $('f, 'v)$ *term* $\Rightarrow$ $('f, 'v)$ *pterm* **where**
to-pterm $(Var\ x) = Var\ x$
$|$ *to-pterm* $(Fun\ f\ ts) = Pfun\ f\ (map\ to\text{-}pterm\ ts)$

Also for contexts.

fun *to-pterm-ctxt* :: $('f, 'v)$ *ctxt* $\Rightarrow$ $('f, 'v)$ *pterm-ctxt* **where**
to-pterm-ctxt $\Box = \Box$
$|$ *to-pterm-ctxt* $(More\ f\ ss1\ C\ ss2) = Cfun\ f\ (map\ to\text{-}pterm\ ss1)\ (to\text{-}pterm\text{-}ctxt\ C)\ (map\ to\text{-}pterm\ ss2)$

2.2 Frequently Used Locales/Contexts

Often certain properties about proof terms only hold when the underlying TRS does not contain variable left-hand sides and/or variables on the right are a subset of the variables on the left and/or the TRS is left-linear.

locale *left-lin* =
fixes R :: $('f, 'v)$ *trs*
assumes *left-lin*:*left-linear-trs* R

locale *no-var-lhs* =
fixes R :: $('f, 'v)$ *trs*
assumes *no-var-lhs*:*Ball* R $(\lambda(l, r).\ is\text{-}Fun\ l)$

locale *var-rhs-subset-lhs* =

fixes R :: $('f,\ 'v)$ *trs*
assumes *varcond*:*Ball* R $(\lambda(l,\ r).$ *vars-term* $r \subseteq$ *vars-term* $l)$

locale *wf-trs* = *no-var-lhs* + *var-rhs-subset-lhs*
locale *left-lin-no-var-lhs* = *left-lin* + *no-var-lhs*
locale *left-lin-wf-trs* = *left-lin* + *wf-trs*

context *wf-trs*
begin
lemma *wf-trs-alt*:
shows *Trs.wf-trs R*
unfolding *wf-trs-def'* **using** *no-var-lhs varcond* **by** *auto*
end

context *left-lin*
begin
lemma *length-var-rule*:
assumes *to-rule* $\alpha \in R$
shows *length* (*var-rule* α) = *length* (*vars-term-list* (*lhs* α))
using *assms*
by (*metis case-prodD left-lin left-linear-trs-def linear-term-var-vars-term-list*)
end

2.3 Proof Term Predicates

The number of arguments of a well-defined proof term over a TRS R using a rule symbol α corresponds to the number of variables in *lhs* α. Also the rewrite rule for α must belong to the TRS R.

inductive-set *wf-pterm* :: $('f,\ 'v)$ *trs* $\Rightarrow$ $('f,\ 'v)$ *pterm set*
for R **where**
[*simp*]: *Var x* $\in$ *wf-pterm R*
|[*intro*]: $\forall$ *t* $\in$ *set ts. t* $\in$ *wf-pterm R* $\Longrightarrow$ *Pfun f ts* $\in$ *wf-pterm R*
|[*intro*]: (*lhs* α, *rhs* α) $\in R \Longrightarrow$ *length As* = *length* (*var-rule* α) $\Longrightarrow$
$\forall$ *a* $\in$ *set As. a* $\in$ *wf-pterm R* $\Longrightarrow$ *Prule* α *As* $\in$ *wf-pterm R*

inductive-set *wf-pterm-ctxt* :: $('f,\ 'v)$ *trs* $\Rightarrow$ $('f,\ 'v)$ *pterm-ctxt set*
for R **where**
[*simp*]: $\Box \in$ *wf-pterm-ctxt R*
|[*intro*]: $\forall$ *s* $\in$ (*set ss1*) $\cup$ (*set ss2*). *s* $\in$ *wf-pterm R* $\Longrightarrow$ *C* $\in$ *wf-pterm-ctxt R* $\Longrightarrow$
Cfun f ss1 C ss2 $\in$ *wf-pterm-ctxt R*
|[*intro*]: (*lhs* α, *rhs* α) $\in R \Longrightarrow$ (*length ss1*) + (*length ss2*) + *1* = *length* (*var-rule* α) $\Longrightarrow$
$\forall$ *s* $\in$ (*set ss1*) $\cup$ (*set ss2*). *s* $\in$ *wf-pterm R* $\Longrightarrow$ *C* $\in$ *wf-pterm-ctxt R* $\Longrightarrow$
Crule α *ss1 C ss2* $\in$ *wf-pterm-ctxt R*

lemma *fun-well-arg*[*intro*]:
assumes *Fun f As* $\in$ *wf-pterm R a* $\in$ *set As*
shows *a* $\in$ *wf-pterm R*
using *assms wf-pterm.cases* **by** *auto*

lemma *trs-well-ctxt-arg*[*intro*]:
  assumes *More f ss1 C ss2* $\in$ *wf-pterm-ctxt R s* $\in$ (*set ss1*) $\cup$ (*set ss2*)
  shows *s* $\in$ *wf-pterm R*
  using *assms wf-pterm-ctxt.cases* **by** *blast*

lemma *trs-well-ctxt-C*[*intro*]:
  assumes *More f ss1 C ss2* $\in$ *wf-pterm-ctxt R*
  shows *C* $\in$ *wf-pterm-ctxt R*
  using *assms wf-pterm-ctxt.cases* **by** *auto*

context *no-var-lhs*
begin
lemma *lhs-is-Fun*:
  assumes *Prule* α *Bs* $\in$ *wf-pterm R*
  shows *is-Fun* (*lhs* α)
 by (*metis Inl-inject assms case-prodD is-FunI is-Prule.simps*(*1*) *is-Prule.simps*(*3*)
is-VarI
   no-var-lhs.no-var-lhs no-var-lhs-axioms term.inject(*2*) *wf-pterm.simps*)
end

lemma *lhs-subst-var-well-def*:
  assumes $\forall$ *i* < *length As. As*!*i* $\in$ *wf-pterm R*
  shows ($\langle As\rangle_\alpha$) *x* $\in$ *wf-pterm R*
proof (*cases map-of* (*zip* (*var-rule* α) *As*) *x*)
  case *None*
  then show *?thesis* **unfolding** *mk-subst-def* **by** *simp*
next
  case (*Some a*)
  then have *a* $\in$ *set As*
   by (*meson in-set-zipE map-of-SomeD*)
  with *assms Some* **show** *?thesis*
   unfolding *mk-subst-def* **using** *in-set-idx* **by** *force*
qed

lemma *lhs-subst-well-def*:
  assumes $\forall$ *i* < *length As. As*!*i* $\in$ *wf-pterm R B* $\in$ *wf-pterm R*
  shows *B* $\cdot$ ($\langle As\rangle_\alpha$) $\in$ *wf-pterm R*
  using *assms* **proof**(*induction B arbitrary*: *As*)
  case (*Var x*)
  then show *?case* **using** *lhs-subst-var-well-def* **by** *simp*
next
  case (*Pfun f Bs*)
  from *Pfun*(*3*) **have** $\forall$ *b* $\in$ *set Bs. b* $\in$ *wf-pterm R*
   by *blast*
  with *Pfun* **show** *?case* **by** *fastforce*
next
  case (*Prule* β *Bs*)
  from *Prule*(*3*) **have** $\forall$ *b* $\in$ *set Bs. b* $\in$ *wf-pterm R*

    by *blast*
  moreover have *length* (*map* ($\lambda t.\ t \cdot \langle As\rangle_\alpha$) *Bs*) = *length* (*var-rule* β)
    using *Prule*(*3*) *wf-pterm.simps* **by** *fastforce*
  moreover from *Prule*(*3*) **have** *to-rule* $\beta \in R$
    using *Inl-inject sum.distinct*(*1*) *wf-pterm.cases* **by** *force*
  ultimately show *?case* **unfolding** *eval-term.simps*(*2*) **using** *Prule*
    by (*simp add*: *wf-pterm.intros*(*3*))
qed

lemma *subt-at-is-wf-pterm*:
  assumes $p \in$ *poss A* **and** $A \in$ *wf-pterm R*
  shows *A*|-*p* $\in$ *wf-pterm R*
  using *assms* **proof**(*induct p arbitrary*:*A*)
  case (*Cons i p*)
  then obtain *f As* **where** *a*:*A* = *Fun f As* **and** *i*:*i* < *length As* **and** *p*:*p* $\in$ *poss* (*As*!*i*)
    using *args-poss* **by** *blast*
  with *Cons*(*3*) **have** *As*!*i* $\in$ *wf-pterm R*
    using *nth-mem* **by** *blast*
  with *Cons.hyps p a* **show** *?case* **by** *simp*
qed *simp*

lemma *ctxt-of-pos-term-well*:
  assumes $p \in$ *poss A* **and** $A \in$ *wf-pterm R*
  shows *ctxt-of-pos-term p A* $\in$ *wf-pterm-ctxt R*
  using *assms* **proof**(*induct p arbitrary*:*A*)
  case (*Cons i p*)
  then obtain *fs As* **where** *a*:*A* = *Fun fs As* **and** *i*:*i* < *length As* **and** *p*:*p* $\in$ *poss* (*As*!*i*)
    using *args-poss* **by** *blast*
  with *Cons*(*3*) **have** *as*:$\forall$ *j* < *length As*. *As*!*j* $\in$ *wf-pterm R*
    using *nth-mem* **by** *blast*
  with *Cons.hyps p i* **have** *IH*:*ctxt-of-pos-term p* (*As*!*i*) $\in$ *wf-pterm-ctxt R*
    by *blast*
  then show *?case* **proof**(*cases fs*)
    case (*Inl* α)
    from *Cons*(*3*) **have** *to-rule* $\alpha \in R$ **unfolding** *a Inl*
      using *wf-pterm.cases* **by** *auto*
    moreover from *Cons*(*3*) *i* **have** *length* (*take i As*) + *length* (*drop* (*Suc i*) *As*) + *1* = *length* (*var-rule* α)
      unfolding *a Inl* **using** *wf-pterm.cases* **by** *force*
    ultimately show *?thesis*
      unfolding *a ctxt-of-pos-term.simps Inl* **using** *as IH wf-pterm-ctxt.intros*(*3*)
    by (*metis* (*no-types*, *opaque-lifting*) *UnE in-set-conv-nth in-set-dropD in-set-takeD*)
  next
    case (*Inr f*)
    show *?thesis*
      unfolding *a ctxt-of-pos-term.simps Inr* **using** *as IH wf-pterm-ctxt.intros*(*2*)
      by (*metis Cons.prems*(*2*) *UnE a fun-well-arg in-set-dropD in-set-takeD*)

qed
qed *simp*

Every normal term is a well-defined proof term.

lemma *to-pterm-wf-pterm*[*simp*]: *to-pterm t* ∈ *wf-pterm R*
by (*induction t*) (*simp-all add*: *wf-pterm.intros*(*2*,*3*))

lemma *to-pterm-trs-ctxt*:
assumes *p* ∈ *poss* (*to-pterm s*)
shows *ctxt-of-pos-term p* (*to-pterm s*) ∈ *wf-pterm-ctxt R*
by (*simp add*: *assms ctxt-of-pos-term-well*)

lemma *to-pterm-ctxt-wf-pterm-ctxt*:
shows *to-pterm-ctxt C* ∈ *wf-pterm-ctxt R*
proof(*induct C*)
case (*More f xs C ys*)
then show *?case* **unfolding** *to-pterm-ctxt.simps*
by (*metis Un-iff fun-well-arg to-pterm.simps*(*2*) *to-pterm-wf-pterm wf-pterm-ctxt.intros*(*2*))
qed *simp*

lemma *ctxt-wf-pterm*:
assumes *A* ∈ *wf-pterm R* **and** *p* ∈ *poss A*
and *B* ∈ *wf-pterm R*
shows (*ctxt-of-pos-term p A*)⟨*B*⟩ ∈ *wf-pterm R*
using *assms* **proof**(*induct p arbitrary*:*A*)
case (*Cons i p*)
from *Cons*(*3*) **obtain** *f As* **where** *A*:*A* = *Fun f As i* < *length As p* ∈ *poss* (*As*!*i*)
using *args-poss* **by** *blast*
moreover with *Cons*(*2*) **have** *As*!*i* ∈ *wf-pterm R*
using *nth-mem* **by** *blast*
ultimately have *IH*:(*ctxt-of-pos-term p* (*As*!*i*))⟨*B*⟩ ∈ *wf-pterm R*
using *Cons.hyps assms*(*3*) **by** *presburger*
from *Cons*(*2*) **have** *as*:∀ *a* ∈ *set As*. *a* ∈ *wf-pterm R*
unfolding *A* **by** *auto*
show *?case* **proof**(*cases f*)
case (*Inl α*)
from *Cons*(*2*) **have** *alpha*:*to-rule α* ∈ *R*
unfolding *A Inl* **using** *wf-pterm.simps* **by** *fastforce*
moreover from *Cons*(*2*) **have** *length As* = *length* (*var-rule α*)
unfolding *A Inl* **using** *wf-pterm.simps* **by** *fastforce*
ultimately show *?thesis*
unfolding *Inl A ctxt-of-pos-term.simps intp-actxt.simps* **using** *wf-pterm.intros*(*3*)[*OF alpha*] *IH as A*(*2*)
by (*smt* (*verit*, *ccfv-SIG*) *id-take-nth-drop in-set-conv-nth le-simps*(*1*) *length-append list.size*(*4*) *nth-append-take nth-append-take-drop-is-nth-conv*)
next
case (*Inr b*)
show *?thesis* **unfolding** *Inr A ctxt-of-pos-term.simps intp-actxt.simps* **using** *wf-pterm.intros*(*2*) *IH as A*(*2*)

by (*smt* (*verit*, *ccfv-SIG*) *Cons-nth-drop-Suc append-take-drop-id in-set-conv-nth length-append length-nth-simps*(*2*) *less-imp-le-nat nth-append-take nth-append-take-drop-is-nth-conv*)
qed
qed *simp*

2.4 'Normal' Terms vs. Proof Terms

lemma *to-pterm-empty*: *is-empty-step* (*to-pterm t*)
proof (*induction t*)
case (*Fun f ts*)
then have *list-all is-empty-step* (*map to-pterm ts*) **using** *list-all-iff* **by** *force*
then show *?case* **by** *simp*
qed *simp*

Variables remain unchanged.

lemma *vars-to-pterm*: *vars-term-list* (*to-pterm t*) = *vars-term-list t*
proof(*induction t*)
case (*Fun f ts*)
then have ∗:*map vars-term-list ts* = *map* (*vars-term-list* ∘ *to-pterm*) *ts* **by** *simp*
show *?case* **by** (*simp add*: ∗ *vars-term-list.simps*(*2*))
qed (*simp add*: *vars-term-list.simps*(*1*))

lemma *poss-list-to-pterm*: *poss-list t* = *poss-list* (*to-pterm t*)
proof(*induction t*)
case (*Fun f ts*)
then have ∗:*map poss-list ts* = *map* (*poss-list* ∘ *to-pterm*) *ts* **by** *simp*
show *?case* **by** (*simp add*: ∗ *poss-list.simps*(*2*))
qed (*simp add*: *poss-list.simps*(*1*))

lemma *p-in-poss-to-pterm*:
assumes $p \in$ *poss t*
shows $p \in$ *poss* (*to-pterm t*)
using *assms poss-list-to-pterm* **by** (*metis poss-list-sound*)

lemma *var-poss-to-pterm*: *var-poss t* = *var-poss* (*to-pterm t*)
proof(*induction t*)
case (*Fun f ts*)
then have ∗:*map var-poss ts* = *map* (*var-poss* ∘ *to-pterm*) *ts* **by** *simp*
then show *?case* **unfolding** *var-poss.simps to-pterm.simps*
by *auto*
qed *simp*

lemma *var-poss-list-to-pterm*: *var-poss-list* (*to-pterm t*) = *var-poss-list t*
proof(*induct t*)
case (*Fun f ts*)
then show *?case* **unfolding** *var-poss-list.simps to-pterm.simps*
by (*metis* (*no-types*, *lifting*) *length-map map-nth-eq-conv nth-mem*)
qed *simp*

to-pterm distributes over application of substitution.

lemma *to-pterm-subst*:
to-pterm $(t \cdot \sigma) = (\textit{to-pterm}\ t) \cdot (\textit{to-pterm} \circ \sigma)$
by (*induct t*, *auto*)

to-pterm distributes over context.

lemma *to-pterm-ctxt-of-pos-apply-term*:
assumes $p \in \textit{poss}\ s$
shows *to-pterm* $((\textit{ctxt-of-pos-term}\ p\ s)\ \langle t\rangle) = (\textit{ctxt-of-pos-term}\ p\ (\textit{to-pterm}\ s))\langle \textit{to-pterm}\ t\rangle$
using *assms* **proof**(*induct p arbitrary*:*s*)
case (*Cons i p*)
then obtain *f ss* **where** s:$s = \textit{Fun}\ f\ ss$ **and** i:$i < \textit{length}\ ss$ **and** p:$p \in \textit{poss}\ (ss!i)$
using *args-poss* **by** *blast*
then show *?case* **unfolding** *s to-pterm.simps ctxt-of-pos-term.simps intp-actxt.simps* **using** *Cons*(*1*)
by (*simp add*: *drop-map take-map*)
qed *simp*

Linear terms become linear proof terms.

lemma *to-pterm-linear*:
assumes *linear-term t*
shows *linear-term* (*to-pterm t*)
using *assms* **proof**(*induction t*)
case (*Fun f ts*)
have $*$:*map vars-term ts* = *map vars-term* (*map to-pterm ts*)
by (*metis* (*mono-tags*, *lifting*) *length-map map-nth-eq-conv set-vars-term-list vars-to-pterm*)
with *Fun* **show** *?case* **by** *auto*
qed *simp*

lemma *lhs-subst-trivial*:
shows *match* $(\textit{to-pterm}\ (\textit{lhs}\ \alpha) \cdot \langle \textit{As}\rangle_\alpha)\ (\textit{to-pterm}\ (\textit{lhs}\ \alpha)) = \textit{Some}\ \langle \textit{As}\rangle_\alpha$
using *match-trivial*
by (*smt comp-def mem-Collect-eq mk-subst-not-mem set-remdups set-rev set-vars-term-list subsetI subst-domain-def vars-to-pterm*)

lemma *to-pterm-ctxt-apply-term*:
to-pterm $C\langle t\rangle = (\textit{to-pterm-ctxt}\ C)\ \langle \textit{to-pterm}\ t\rangle$
by(*induct C*) *simp-all*

2.5 Substitutions

lemma *fun-mk-subst*[*simp*]:
assumes $\forall x.\ f\ (\textit{Var}\ x) = \textit{Var}\ x$
shows $f \circ (\textit{mk-subst Var}\ (\textit{zip vs ts})) = \textit{mk-subst Var}\ (\textit{zip vs}\ (\textit{map}\ f\ ts))$
proof−
have $\forall a.\ f\ (\textit{case map-of}\ (\textit{zip vs ts})\ a\ \textit{of None} \Rightarrow \textit{Var}\ a \mid \textit{Some}\ t \Rightarrow t)$
$= (\textit{case map-of}\ (\textit{zip vs ts})\ a\ \textit{of None} \Rightarrow \textit{Var}\ a \mid \textit{Some}\ t \Rightarrow f\ t)$
using *assms* **by** (*simp add*: *option.case-eq-if*)

moreover have $\forall a.$ (*case map-of* (*zip vs* (*map f ts*)) *a of None* $\Rightarrow$ *Var a* | *Some x* $\Rightarrow$ *x*)
= (*case* (*map-of* (*zip vs ts*)) *a of None* $\Rightarrow$ *Var a* | *Some t* $\Rightarrow$ *f t*)
by (*simp add*:*zip-map2 map-of-map option.case-eq-if option.map-sel*)
ultimately show *?thesis* **unfolding** *mk-subst-def* **unfolding** *comp-def* **by** *auto*
qed

lemma *apply-lhs-subst-var-rule*:
assumes *length ts* = *length* (*var-rule* α)
shows *map* ($\langle ts\rangle_\alpha$) (*var-rule* α) = *ts*
using *assms* **by** (*simp add*: *mk-subst-distinct map-nth-eq-conv*)

lemma *match-lhs-subst*:
assumes *match B* (*to-pterm* (*lhs* α)) = *Some* σ
shows $\exists$ *Bs*. *length Bs* = *length* (*var-rule* α) $\wedge$
B = (*to-pterm* (*lhs* α)) $\cdot$ $\langle Bs\rangle_\alpha$ $\wedge$
($\forall x \in$ *set* (*var-rule* α). σ *x* = ($\langle Bs\rangle_\alpha$) *x*)
proof−
obtain *Bs* **where** *Bs*:*length Bs* = *length* (*var-rule* α)
$\forall i <$ *length* (*var-rule* α). *Bs*!*i* = σ ((*var-rule* α)!*i*)
using *length-map nth-map* **by** *blast*
then have *2*:($\forall x \in$ *set* (*var-rule* α). σ *x* = ($\langle Bs\rangle_\alpha$) *x*)
by (*smt apply-lhs-subst-var-rule in-set-idx nth-map*)
have *v*:*vars-term* (*to-pterm* (*lhs* α)) = *set* (*var-rule* α)
by (*metis comp-apply set-remdups set-rev set-vars-term-list vars-to-pterm*)
from *assms* **have** *B* = (*to-pterm* (*lhs* α)) $\cdot$ σ
using *match-matches* **by** *blast*
also have ... = (*to-pterm* (*lhs* α)) $\cdot$ $\langle Bs\rangle_\alpha$
by (*intro term-subst-eq*, *insert 2 v*, *auto*)
finally show *?thesis* **using** *Bs 2* **by** *auto*
qed

lemma *apply-subst-wf-pterm*:
assumes *A* $\in$ *wf-pterm R*
and $\forall x \in$ *vars-term A*. σ *x* $\in$ *wf-pterm R*
shows *A* $\cdot$ σ $\in$ *wf-pterm R*
using *assms* **proof**(*induct A*)
case (*2 ts f*)
{fix *t* **assume** *t*:*t* $\in$ *set ts*
with *2*(*2*) **have** ($\forall x{\in}$*vars-term t*. σ *x* $\in$ *wf-pterm R*)
by (*meson term.set-intros*(*4*))
with *t 2*(*1*) **have** *t* $\cdot$ σ $\in$ *wf-pterm R*
by *blast*
}
then show *?case* **unfolding** *eval-term.simps* **by** (*simp add*: *wf-pterm.intros*(*2*))
next
case (*3* α *As*)
{fix *a* **assume** *a*:*a* $\in$ *set As*
with *3*(*4*) **have** ($\forall x{\in}$*vars-term a*. σ *x* $\in$ *wf-pterm R*)

by (*meson term.set-intros*(*4*))
with *a 3*(*3*) **have** $a \cdot \sigma \in$ *wf-pterm R*
by *blast*
}
with *3*(*1*,*2*) **show** *?case* **unfolding** *eval-term.simps* **by** (*simp add*: *wf-pterm.intros*(*3*))
qed *simp*

lemma *subst-well-def*:
assumes $B \in$ *wf-pterm R* $A \cdot \sigma = B$ $x \in$ *vars-term A*
shows $\sigma\ x \in$ *wf-pterm R*
using *assms* **by** (*metis* (*no-types*, *lifting*) *poss-imp-subst-poss eval-term.simps*(*1*) *subt-at-is-wf-pterm subt-at-subst vars-term-poss-subt-at*)

lemma *lhs-subst-args-wf-pterm*:
assumes *to-pterm* (*lhs* α) $\cdot \langle As\rangle_\alpha \in$ *wf-pterm R* **and** *length As* = *length* (*var-rule* α)
shows $\forall\, a \in$ *set As*. $a \in$ *wf-pterm R*
proof−
from *assms* **have** *map* ($\langle As\rangle_\alpha$) (*var-rule* α) = *As*
using *apply-lhs-subst-var-rule* **by** *blast*
with *assms* **show** *?thesis*
by (*smt comp-apply in-set-idx map-nth-eq-conv nth-mem set-remdups set-rev set-vars-term-list subst-well-def vars-to-pterm*)
qed

lemma *match-well-def*:
assumes $B \in$ *wf-pterm R match B A* = *Some* σ
shows $\forall\, i <$ *length* (*vars-distinct A*). σ ((*vars-distinct A*) ! *i*) $\in$ *wf-pterm R*
using *assms subst-well-def match-matches*
by (*smt comp-apply nth-mem set-remdups set-rev set-vars-term-list*)

lemma *subst-imp-well-def*:
assumes $A \cdot \sigma \in$ *wf-pterm R*
shows $A \in$ *wf-pterm R*
using *assms* **proof**(*induct A*)
case (*Pfun f As*)
{fix *i* **assume** *i*:$i <$ *length As*
with *Pfun*(*2*) **have** (*As*!*i*) $\cdot \sigma \in$ *wf-pterm R*
by *auto*
with *Pfun*(*1*) *i* **have** *As*!*i* $\in$ *wf-pterm R*
by *simp*
}
then show *?case* **using** *wf-pterm.intros*(*2*)
by (*metis in-set-idx*)
next
case (*Prule* α *As*)
{fix *i* **assume** *i*:$i <$ *length As*
with *Prule*(*2*) **have** (*As*!*i*) $\cdot \sigma \in$ *wf-pterm R*
by *auto*

with *Prule*(*1*) *i* **have** *As*!*i* ∈ *wf-pterm R*
by *simp*
}
moreover from *Prule*(*2*) **have** *to-rule* α ∈ *R length As* = *length* (*var-rule* α)
using *wf-pterm.cases* **by** *force*+
ultimately show *?case* **using** *wf-pterm.intros*(*3*) *Prule*(*2*)
by (*metis in-set-idx*)
qed *simp*

lemma *lhs-subst-var-i*:
assumes *x* = (*var-rule* α)!*i* **and** *i* < *length* (*var-rule* α) **and** *i* < *length As*
shows ($\langle As\rangle_\alpha$) *x* = *As*!*i*
using *assms mk-subst-distinct distinct-remdups* **by** (*metis comp-apply distinct-rev*)

lemma *lhs-subst-not-var-i*:
assumes ¬(∃ *i* < *length As*. *i* < *length* (*var-rule* α) ∧ *x* = (*var-rule* α)!*i*)
shows ($\langle As\rangle_\alpha$) *x* = *Var x*
using *assms* **proof**(*rule contrapos-np*)
{assume ($\langle As\rangle_\alpha$) *x* ≠ *Var x*
then obtain *i* **where** *i* < *length* (*zip* (*var-rule* α) *As*) **and** (*var-rule* α)!*i* = *x*
unfolding *mk-subst-def* **by** (*smt assms imageE in-set-zip map-of-eq-None-iff option.case-eq-if*)
then show ∃ *i*<*length As*. *i* < *length* (*var-rule* α) ∧ *x* = *var-rule* α ! *i*
by *auto*
}
qed

lemma *lhs-subst-upd*:
assumes *length ss1* < *length* (*var-rule* α)
shows (($\langle ss1$ @ *t* # *ss2*$\rangle_\alpha$) ((*var-rule* α)!(*length ss1*) := *s*)) = $\langle ss1$ @ *s* # *ss2*$\rangle_\alpha$
proof
fix *x*
show (($\langle ss1$ @ *t* # *ss2*$\rangle_\alpha$)(*var-rule* α ! *length ss1* := *s*)) *x* = ($\langle ss1$ @ *s* # *ss2*$\rangle_\alpha$) *x* **proof**(*cases x* = (*var-rule* α)!(*length ss1*))
case *True*
with *assms* **have** (($\langle ss1$ @ *t* # *ss2*$\rangle_\alpha$)(*var-rule* α ! *length ss1* := *s*)) *x* = *s*
by *simp*
moreover from *assms* **have** ($\langle ss1$ @ *s* # *ss2*$\rangle_\alpha$) *x* = *s* **unfolding** *True*
by (*smt* (*verit*, *del-insts*) *add.commute add-Suc-right le-add-same-cancel2 le-imp-less-Suc length-append length-nth-simps*(*2*) *lhs-subst-var-i nth-append-length zero-order*(*1*))
ultimately show *?thesis* **by** *simp*
next
case *False*
then show *?thesis*
by (*smt* (*verit*, *del-insts*) *append-Cons-nth-not-middle fun-upd-apply length-append length-nth-simps*(*2*) *lhs-subst-not-var-i lhs-subst-var-i*)
qed
qed

lemma *eval-lhs-subst*:
assumes *l*:*length* (*var-rule* α) = *length As*
shows (*to-pterm* (*lhs* α)) $\cdot$ $\langle As\rangle_\alpha$ $\cdot$ σ = (*to-pterm* (*lhs* α)) $\cdot$ $\langle$*map* ($\lambda a.\ a \cdot \sigma$) $As\rangle_\alpha$
proof−
{**fix** x **assume** $x \in$ *vars-term* (*to-pterm* (*lhs* α))
then obtain i **where** i:$i <$ *length* (*var-rule* α) (*var-rule* α) !$i = x$
using *vars-to-pterm* **by** (*metis in-set-conv-nth set-vars-term-list vars-term-list-vars-distinct*)

with l **have** ($\langle As\rangle_\alpha$) x = *As*!i
by (*metis lhs-subst-var-i*)
then have *1*:($\langle As\rangle_\alpha \circ_s \sigma$) x = *As*!$i \cdot \sigma$
unfolding *subst-compose-def* **by** *simp*
from i l **have** ($\langle$*map* ($\lambda a.\ a \cdot \sigma$) $As\rangle_\alpha$) x = *map* ($\lambda a.\ a \cdot \sigma$) *As* ! i
using *lhs-subst-var-i* **by** (*metis length-map*)
with *1* i l **have** ($\langle As\rangle_\alpha \circ_s \sigma$) x = ($\langle$*map* ($\lambda a.\ a \cdot \sigma$) $As\rangle_\alpha$) x **by** *simp*
}
then show *?thesis*
by (*smt* (*verit*, *ccfv-SIG*) *eval-same-vars-cong subst-subst-compose*)
qed

lemma *var-rule-pos-subst*:
assumes $i <$ *length* (*var-rule* α) *length ss* = *length* (*var-rule* α)
and $p \in$ *poss* (*lhs* α) *Var* ((*var-rule* α)!i) = (*lhs* α)|-p
shows *lhs* $\alpha \cdot \langle ss\rangle_\alpha$ |- (p@q) = (*ss*!i)|-q
proof−
from *assms*(*1*,*2*) **have** ($\langle ss\rangle_\alpha$) ((*var-rule* α)!i) = *ss*!i
using *lhs-subst-var-i* **by** *force*
with *assms*(*3*,*4*) **show** *?thesis* **by** *auto*
qed

lemma *lhs-subst-var-rule*:
assumes *vars-term* $t \subseteq$ *vars-term* (*lhs* α)
shows $t \cdot \langle$*map* σ (*var-rule* α)$\rangle_\alpha$ = $t \cdot \sigma$
using *assms* **by** (*smt* (*verit*, *ccfv-SIG*) *apply-lhs-subst-var-rule comp-apply length-map map-eq-conv set-remdups set-rev set-vars-term-list subsetD term-subst-eq-conv*)

2.6 Contexts

lemma *match-lhs-context*:
assumes $i <$ *length* (*vars-term-list* t) $\wedge$ p = (*var-poss-list* t)!i
and *linear-term* t
and *match* (((*ctxt-of-pos-term* p ($t \cdot \sigma$)))$\langle B\rangle$) t = *Some* τ
shows *map* τ (*vars-term-list* t) = (*map* σ (*vars-term-list* t))[i := B]
proof−
from *assms* **have** (*ctxt-of-pos-term* p ($t \cdot \sigma$))$\langle B\rangle$ = $t \cdot$ (σ(*vars-term-list* t!i := B))
using *ctxt-apply-term-subst* **by** *blast*

with *assms*(*3*) **have** ∗:(∀ *x*∈*vars-term t*. (σ(*vars-term-list t*!*i* := *B*)) *x* = τ *x*)
using *match-complete′* **by** (*metis option.inject*)
from *assms*(*2*) **have** *distinct* (*vars-term-list t*)
by (*metis distinct-remdups distinct-rev linear-term-var-vars-term-list o-apply*)
with ∗ *assms*(*1*) **show** *?thesis*
by (*smt* (*verit*, *ccfv-threshold*) *fun-upd-other fun-upd-same length-list-update length-map map-nth-eq-conv nth-eq-iff-index-eq nth-list-update nth-mem set-vars-term-list*)
qed

lemma *ctxt-lhs-subst*:
assumes *i*:*i* < *length* (*var-poss-list* (*lhs* α)) **and** *l*:*length As* = *length* (*var-rule* α)
and *p1*:*p1* = *var-poss-list* (*lhs* α) ! *i* **and** *lin*:*linear-term* (*lhs* α)
and *p2* ∈ *poss* (*As*!*i*)
shows (*ctxt-of-pos-term* (*p1* @ *p2*) (*to-pterm* (*lhs* α) · ⟨*As*⟩$_\alpha$))⟨*A*⟩ = (*to-pterm* (*lhs* α)) · ⟨*take i As* @ (*ctxt-of-pos-term p2* (*As*!*i*))⟨*A*⟩ # *drop* (*Suc i*) *As*⟩$_\alpha$
proof−
have *l2*:*length* (*var-poss-list* (*lhs* α)) = *length* (*var-rule* α)
using *lin* **by** (*metis length-var-poss-list linear-term-var-vars-term-list*)
from *p1 i* **have** *p1-pos*:*p1* ∈ *poss* (*to-pterm* (*lhs* α))
by (*metis nth-mem var-poss-imp-poss var-poss-list-sound var-poss-to-pterm*)
have *sub*:(*to-pterm* (*lhs* α))|-*p1* = *Var* (*vars-term-list* (*lhs* α)!*i*)
by (*metis i length-var-poss-list p1 var-poss-list-to-pterm vars-term-list-var-poss-list vars-to-pterm*)
have ∗∗: (*to-pterm* (*lhs* α) · ⟨*As*⟩$_\alpha$)|-*p1* = *As*!*i*
unfolding *subt-at-subst*[*OF p1-pos*] *sub eval-term.simps* **using** *i l l2* **by** (*metis lhs-subst-var-i lin linear-term-var-vars-term-list*)
then have ∗:(*ctxt-of-pos-term* (*p1* @ *p2*) (*to-pterm* (*lhs* α) · ⟨*As*⟩$_\alpha$)) = ((*ctxt-of-pos-term p1* (*to-pterm* (*lhs* α))) ·$_c$ ⟨*As*⟩$_\alpha$) ∘$_c$ (*ctxt-of-pos-term p2* (*As*!*i*))
using *ctxt-of-pos-term-append ctxt-of-pos-term-subst* **by** (*metis p1-pos poss-imp-subst-poss*)

show *?thesis*
by (*smt* (*verit*, *ccfv-threshold*) ∗ ∗∗ *ctxt-ctxt-compose*
ctxt-subst-apply lhs-subst-upd append-Cons-nth-not-middle
i id-take-nth-drop l l2 less-imp-le-nat lin linear-term-var-vars-term-list
nth-append-take p1-pos sub to-pterm-linear
ctxt-of-pos-term-append ctxt-supt-id eval-term.simps(*1*) *poss-imp-subst-poss*
replace-at-append-subst subt-at-subst)
qed

lemma *ctxt-rule-obtain-pos*:
assumes *iq*:*i*#*q* ∈ *poss* (*Prule* α *As*)
and *p-pos*:*p* ∈ *poss* (*source* (*Prule* α *As*))
and *ctxt*:*source-ctxt* (*ctxt-of-pos-term* (*i*#*q*) (*Prule* α *As*)) = *ctxt-of-pos-term p* (*source* (*Prule* α *As*))
and *lin*:*linear-term* (*lhs* α)
and *l*:*length As* = *length* (*var-rule* α)
shows ∃ *p1 p2*. *p* = *p1*@*p2* ∧ *p1* = *var-poss-list* (*lhs* α)!*i* ∧ *p2* ∈ *poss* (*source*

$(As!i))$
proof −
  from *iq* **have** $i{:}i < \mathit{length}\ As$
    by *simp*
  let $?p1{=}\mathit{var\text{-}poss\text{-}list}\ (\mathit{lhs}\ \alpha)!i$
  have $p1{:}(\mathit{var\text{-}poss\text{-}list}\ (\mathit{lhs}\ \alpha)\ !\ \mathit{length}\ (\mathit{take}\ i\ As)) = ?p1$
    using *i* **by** *fastforce*
  have $p1\text{-}pos{:}?p1 \in \mathit{poss}\ (\mathit{lhs}\ \alpha)$
    by (*metis i l length-var-poss-list lin linear-term-var-vars-term-list nth-mem var-poss-imp-poss var-poss-list-sound*)
  then have $*{:}\mathit{source\text{-}ctxt}\ (\mathit{ctxt\text{-}of\text{-}pos\text{-}term}\ (i\ \#\ q)\ (\mathit{Prule}\ \alpha\ As)) = ((\mathit{ctxt\text{-}of\text{-}pos\text{-}term}\ ?p1\ (\mathit{lhs}\ \alpha)) \cdot_c \langle \mathit{map}\ \mathit{source}\ (\mathit{take}\ i\ As\ @\ \mathit{Var}\ (\mathit{vars\text{-}term\text{-}list}\ (\mathit{lhs}\ \alpha)\ !\ \mathit{length}\ (\mathit{take}\ i\ As))\ \#\ \mathit{drop}\ (\mathit{Suc}\ i)\ As)\rangle_\alpha) \circ_c$
    $\mathit{source\text{-}ctxt}\ (\mathit{ctxt\text{-}of\text{-}pos\text{-}term}\ q\ (As\ !\ i))$
    unfolding *ctxt-of-pos-term.simps source-ctxt.simps Let-def p1* **by** (*simp add*: *ctxt-of-pos-term-subst*)
  from *ctxt* **have** $?p1\ \leq_p\ p$
    unfolding $*$ **using** *p1-pos p-pos* **unfolding** *source.simps* **using** *ctxt-subst-comp-pos* **by** *blast*
  then obtain *p2* **where** $p{:}p = ?p1@p2$
    using *less-eq-pos-def* **by** *force*
  have $(\mathit{lhs}\ \alpha)|\text{-}?p1 = \mathit{Var}\ (\mathit{vars\text{-}term\text{-}list}\ (\mathit{lhs}\ \alpha)\ !i)$
    by (*metis i l lin linear-term-var-vars-term-list vars-term-list-var-poss-list*)
  moreover have $\mathit{Var}\ (\mathit{vars\text{-}term\text{-}list}\ (\mathit{lhs}\ \alpha)\ !i) \cdot \langle \mathit{map}\ \mathit{source}\ As\rangle_\alpha = \mathit{source}\ (As!i)$
    unfolding *eval-term.simps* **using** *lhs-subst-var-i i l* **by** (*smt* (*verit*, *best*) *length-map lin linear-term-var-vars-term-list nth-map*)
  ultimately have $p2 \in \mathit{poss}\ (\mathit{source}\ (As!i))$
    using *p-pos* **unfolding** *p* **using** *p1-pos* **by** *auto*
  with *p* **show** *?thesis* **by** *simp*
qed

2.7 Source and Target

lemma *source-empty-step*:
  assumes *is-empty-step t*
  shows $\mathit{to\text{-}pterm}\ (\mathit{source}\ t) = t$
using *assms* **by** (*induction t*) (*simp-all add*: *list-all-length map-nth-eq-conv*)

lemma *empty-coinitial*:
  shows $\mathit{co\text{-}initial}\ A\ t \Longrightarrow \mathit{is\text{-}empty\text{-}step}\ t \Longrightarrow \mathit{to\text{-}pterm}\ (\mathit{source}\ A) = t$
  by (*simp add*: *source-empty-step*)

lemma *source-to-pterm*[*simp*]: $\mathit{source}\ (\mathit{to\text{-}pterm}\ t) = t$
  by (*induction t*) (*simp-all add*: *map-nth-eq-conv*)

lemma *target-to-pterm*[*simp*]: $\mathit{target}\ (\mathit{to\text{-}pterm}\ t) = t$
  by (*induction t*) (*simp-all add*: *map-nth-eq-conv*)

```
lemma vars-term-source:
  assumes A ∈ wf-pterm R
  shows vars-term A = vars-term (source A)
  using assms proof(induct A)
  case (3 α As)
  show ?case proof
    {fix x assume x ∈ vars-term (Prule α As)
      then obtain i where i:i < length As x ∈ vars-term (As!i)
        by (metis term.sel(4) var-imp-var-of-arg)
    from i(1) 3(2) obtain j where j:j < length (vars-term-list (lhs α)) vars-term-list
(lhs α)!j = var-rule α !i
        by (metis comp-apply in-set-idx nth-mem set-remdups set-rev)
      let ?p=(var-poss-list (lhs α)!j)
      from j have p:?p ∈ poss (lhs α)
      by (metis in-set-conv-nth length-var-poss-list var-poss-imp-poss var-poss-list-sound)

      with 3(2) i(1) j have source (Prule α As) |- ?p = source (As!i)
        using mk-subst-distinct unfolding source.simps
          by (smt (verit, best) comp-apply distinct-remdups distinct-rev filter-cong
length-map map-nth-conv mk-subst-same eval-term.simps(1) subt-at-subst vars-term-list-var-poss-list)

      with 3(3) have x ∈ vars-term (source (Prule α As))
        unfolding source.simps using vars-term-subt-at p
        by (smt (verit, ccfv-SIG) i nth-mem poss-imp-subst-poss subsetD)
    }
    then show vars-term (Prule α As) ⊆ vars-term (source (Prule α As))
      by blast
    {fix x assume x ∈ vars-term (source (Prule α As))
     then obtain y where y:y ∈ vars-term (lhs α)  x ∈ vars-term (((⟨map source
As⟩_α) y)
        using vars-term-subst by force
      then obtain i where i:i < length (var-rule α) y = var-rule α!i
        by (metis in-set-idx set-vars-term-list vars-term-list-vars-distinct)
      with y(2) 3(2) have x ∈ vars-term (source (As!i))
        by (simp add: mk-subst-distinct)
      with 3 i(1) have x ∈ vars-term (Prule α As)
        by (metis nth-mem term.set-intros(4))
    }
    then show vars-term (source (Prule α As)) ⊆ vars-term (Prule α As)
      by blast
  qed
qed auto

context var-rhs-subset-lhs
begin
lemma vars-term-target:
  assumes A ∈ wf-pterm R
  shows vars-term (target A) ⊆ vars-term A
  using assms proof(induct A)
```

case (*3 α As*)
show *?case* **proof**
fix *x* **assume** *x* ∈ *vars-term* (*target* (*Prule α As*))
then obtain *y* **where** *y*:*y* ∈ *vars-term* (*rhs α*) *x* ∈ *vars-term* (((⟨*map target As*⟩$_{\alpha}$) *y*)
using *vars-term-subst* **by** *force*
then have *y* ∈ *vars-term* (*lhs α*)
using *3.hyps*(*1*) *varcond* **by** *auto*
then obtain *i* **where** *i*:*i* < *length* (*var-rule α*) *y* = *var-rule α*!*i*
by (*metis in-set-idx set-vars-term-list vars-term-list-vars-distinct*)
with *y*(*2*) *3*(*2*) **have** *x* ∈ *vars-term* (*target* (*As*!*i*))
by (*simp add*: *mk-subst-distinct*)
with *3 i*(*1*) **show** *x* ∈ *vars-term* (*Prule α As*)
by *fastforce*
qed
qed *auto*
end

lemma *linear-source-imp-linear-pterm*:
assumes *A* ∈ *wf-pterm R linear-term* (*source A*)
shows *linear-term A*
using *assms* **proof**(*induct A*)
case (*2 As f*)
then show *?case* **unfolding** *source.simps linear-term.simps* **using** *vars-term-source*
by (*smt* (*verit*, *ccfv-SIG*) *in-set-idx length-map map-equality-iff nth-map nth-mem*)
next
case (*3 α As*)
{**fix** *a* **assume** *a*:*a* ∈ *set As*
with *3*(*2*) **obtain** *i* **where** *i*:*i* < *length* (*var-rule α*) *As*!*i* = *a*
by (*metis in-set-idx*)
let *?x*=*var-rule α* ! *i*
from *i* **have** *?x* ∈ *vars-term* (*lhs α*)
by (*metis comp-apply nth-mem set-remdups set-rev set-vars-term-list*)
then obtain *p* **where** *p* ∈ *poss* (*lhs α*) *lhs α* |-*p* = *Var ?x*
by (*meson vars-term-poss-subt-at*)
then have *source* (*Prule α As*) ⊵ *source a*
unfolding *source.simps* **using** *lhs-subst-var-i*[*of ?x α i As*] *i 3*(*2*)
by (*smt* (*verit*, *best*) ‹*var-rule α* ! *i* ∈ *vars-term* (*lhs α*)› *apply-lhs-subst-var-rule eval-term.simps*(*1*) *length-map map-nth-conv supteq-subst vars-term-supteq*)
then have *linear-term* (*source a*)
using *3*(*4*) **by** (*metis subt-at-linear supteq-imp-subt-at*)
with *3*(*3*) *a* **have** *linear-term a* **by** *simp*
}
moreover have *is-partition* (*map vars-term As*) **proof**−
{**fix** *i j* **assume** *i*:*i* < *length As* **and** *j*:*j* < *length As* **and** *ij*:*i* ≠ *j*
let *?x*=*var-rule α* ! *i* **and** *?y*=*var-rule α* ! *j*
from *i j ij 3*(*2*) **have** *xy*:*?x* ≠ *?y*
by (*simp add*: *nth-eq-iff-index-eq*)
from *i 3*(*2*) **have** *?x* ∈ *vars-term* (*lhs α*)

by (*metis comp-apply nth-mem set-remdups set-rev set-vars-term-list*)
then obtain *p* **where** *p*:$p \in$ *poss* (*lhs* α) *lhs* α |-*p* = *Var ?x*
by (*meson vars-term-poss-subt-at*)
from *j 3*(*2*) **have** *?y* $\in$ *vars-term* (*lhs* α)
by (*metis comp-apply nth-mem set-remdups set-rev set-vars-term-list*)
then obtain *q* **where** *q*:$q \in$ *poss* (*lhs* α) *lhs* α |-*q* = *Var ?y*
by (*meson vars-term-poss-subt-at*)
from *xy p q* **have** $p \perp q$
using *less-eq-pos-def parallel-pos* **by** *auto*
moreover have *source* (*Prule* α *As*) |-*p* = *source* (*As*!*i*)
unfolding *source.simps* **by** (*metis* (*mono-tags*, *lifting*) *3.hyps*(*2*) *eval-term.simps*(*1*) *i length-map lhs-subst-var-i nth-map p subt-at-subst*)
moreover have *source* (*Prule* α *As*) |-*q* = *source* (*As*!*j*)
unfolding *source.simps* **by** (*metis* (*mono-tags*, *lifting*) *3.hyps*(*2*) *eval-term.simps*(*1*) *j length-map lhs-subst-var-i nth-map q subt-at-subst*)
ultimately have *vars-term* (*source* (*As*!*i*)) $\cap$ *vars-term* (*source* (*As*!*j*)) = {}
using *3*(*4*) **by** (*metis linear-subterms-disjoint-vars p*(*1*) *poss-imp-subst-poss q*(*1*) *source.simps*(*3*))
then have *vars-term* (*As*!*i*) $\cap$ *vars-term* (*As*!*j*) = {}
using *vars-term-source 3*(*3*) *i j* **using** *nth-mem* **by** *blast*
}
then show *?thesis*
unfolding *is-partition-alt is-partition-alt-def* **by** *simp*
qed
ultimately show *?case* **unfolding** *source.simps linear-term.simps* **by** *simp*
qed *simp*

context *var-rhs-subset-lhs*
begin
lemma *target-apply-subst*:
assumes $A \in$ *wf-pterm R*
shows *target* ($A \cdot \sigma$) = (*target A*) $\cdot$ (*target* $\circ\ \sigma$)
using *assms*(*1*) **proof**(*induct A*)
case (*2 ts f*)
then have (*map target* (*map* ($\lambda t.\ t \cdot \sigma$) *ts*)) = (*map* ($\lambda t.\ t \cdot$ (*target* $\circ\ \sigma$)) (*map target ts*))
unfolding *map-map o-def* **by** *auto*
then show *?case* **unfolding** *eval-term.simps target.simps* **by** *simp*
next
case (*3* α *As*)
have *id*:$\forall\ x \in$ *vars-term* (*rhs* α). ($\langle$*map* (*target* $\circ$ ($\lambda t.\ t \cdot \sigma$)) *As*$\rangle_\alpha$) *x* = ($\langle$*map target As*$\rangle_\alpha \circ_s$ (*target* $\circ\ \sigma$)) *x*
proof−
have *vars*:*vars-term* (*rhs* α) $\subseteq$ *set* (*var-rule* α)
using *3*(*1*) *varcond* **by** *auto*
{ fix *i* **assume** *i*:$i <$ *length* (*var-rule* α)
with *3* **have** ($\langle$*map* (*target* $\circ$ ($\lambda t.\ t \cdot \sigma$)) *As*$\rangle_\alpha$) ((*var-rule* α)!*i*) = *target* ((*As*!*i*) $\cdot\ \sigma$)
by (*simp add*: *mk-subst-distinct*)

also have ... = $target\ (As!i) \cdot (target \circ \sigma)$
using *3 i* **by** (*metis nth-mem*)
also have ... = $(\langle map\ target\ As\rangle_\alpha \circ_s (target \circ \sigma))\ ((var\text{-}rule\ \alpha)!i)$
using *3 i* **unfolding** *subst-compose-def* **by** (*simp add*: *mk-subst-distinct*)
finally have $(\langle map\ (target \circ (\lambda t.\ t \cdot \sigma))\ As\rangle_\alpha)\ ((var\text{-}rule\ \alpha)!i) = (\langle map\ target\ As\rangle_\alpha \circ_s (target \circ \sigma))\ ((var\text{-}rule\ \alpha)!i)$ **.**
} **with** *vars* **show** *?thesis* **by** (*smt* (*z3*) *in-mono in-set-conv-nth*)
qed
have $target\ ((Prule\ \alpha\ As) \cdot \sigma) = (rhs\ \alpha) \cdot \langle map\ (target \circ (\lambda t.\ t \cdot \sigma))\ As\rangle_\alpha$
unfolding *eval-term.simps*(*2*) **by** *simp*
also have ... = $(rhs\ \alpha) \cdot (\langle map\ target\ As\rangle_\alpha \circ_s (target \circ \sigma))$
using *id* **by** (*meson term-subst-eq*)
also have ... = $(target\ (Prule\ \alpha\ As)) \cdot (target \circ \sigma)$ **by** *simp*
finally show *?case* **.**
qed *simp*
end

context *var-rhs-subset-lhs*
begin
lemma *tgt-subst-simp*:
assumes $A \in wf\text{-}pterm\ R$
shows $target\ (A \cdot \sigma) = target\ ((to\text{-}pterm\ (target\ A)) \cdot \sigma)$
by (*metis assms target-apply-subst target-to-pterm to-pterm-wf-pterm*)
end

lemma *target-empty-apply-subst*:
assumes *is-empty-step t*
shows $target\ (t \cdot \sigma) = (target\ t) \cdot (target \circ \sigma)$
using *assms* **proof**(*induction t*)
case (*Var x*)
then show *?case* **by** (*metis comp-apply eval-term.simps*(*1*) *target.simps*(*1*))
next
case (*Pfun f As*)
from *Pfun*(*2*) **have** $\forall a \in set\ As.\ is\text{-}empty\text{-}step\ a$
by (*simp add*: *Ball-set-list-all*)
with *Pfun*(*1*) **show** *?case* **by** *simp*
next
case (*Prule* α *As*)
then show *?case*
using *is-empty-step.simps*(*3*) **by** *blast*
qed

lemma *source-ctxt-comp*:$source\text{-}ctxt\ (C1 \circ_c C2) = source\text{-}ctxt\ C1 \circ_c source\text{-}ctxt\ C2$
by(*induct C1*) (*simp-all add*:*ctxt-monoid-mult.mult-assoc*)

lemma *context-source*: $source\ (A\langle B\rangle) = source\ (A\langle to\text{-}pterm\ (source\ B)\rangle)$
proof(*induct A rule*:*actxt.induct*)
case (*More f ss1 A ss2*)

```
  then show ?case by(cases f) simp-all
qed simp

lemma context-target: target (A⟨B⟩) = target (A⟨to-pterm (target B)⟩)
proof(induct A rule:actxt.induct)
  case (More f ss1 A ss2)
  then show ?case by(cases f) simp-all
qed simp

lemma source-to-pterm-ctxt:
  source ((to-pterm-ctxt C)⟨A⟩) = C⟨source A⟩
  by (metis context-source source-to-pterm to-pterm-ctxt-apply-term)

lemma target-to-pterm-ctxt:
  target ((to-pterm-ctxt C)⟨A⟩) = C⟨target A⟩
  by (metis context-target target-to-pterm to-pterm-ctxt-apply-term)

lemma source-ctxt-to-pterm:
  assumes p ∈ poss s
  shows source-ctxt (ctxt-of-pos-term p (to-pterm s)) = ctxt-of-pos-term p s
using assms proof(induct p arbitrary:s)
  case (Cons i p)
  then obtain f ss where s:s = Fun f ss and i < length ss and p ∈ poss (ss!i)
    using args-poss by blast
  then show ?case
     unfolding s to-pterm.simps ctxt-of-pos-term.simps source-ctxt.simps using
Cons(1)
   by (smt (verit, best) drop-map nth-map source.simps(2) source-to-pterm take-map
term.inject(2) to-pterm.simps(2))
qed simp

lemma to-pterm-ctxt-at-pos:
  assumes p ∈ poss s
  shows  ctxt-of-pos-term p (to-pterm s) = to-pterm-ctxt (ctxt-of-pos-term p s)
using assms proof(induct p arbitrary:s)
  case (Cons i p)
  then obtain f ss where s:s = Fun f ss
    using args-poss by blast
  with Cons show ?case
    using drop-map s take-map by force
qed simp

lemma to-pterm-ctxt-hole-pos: hole-pos C = hole-pos (to-pterm-ctxt C)
  by(induct C) simp-all

lemma source-to-pterm-ctxt':
  assumes q ∈ poss s
  shows source-ctxt (to-pterm-ctxt (ctxt-of-pos-term q s)) = ctxt-of-pos-term q s
using assms proof(induct q arbitrary: s)
```

case $(Cons\ i\ q)$
then obtain $f\ ss$ **where** $s{:}s = Fun\ f\ ss$ **and** $i{:}i < length\ ss$
by $(meson\ args\text{-}poss)$
with $Cons$ **have** $IH{:}source\text{-}ctxt\ (to\text{-}pterm\text{-}ctxt\ (ctxt\text{-}of\text{-}pos\text{-}term\ q\ (ss!i))) = ctxt\text{-}of\text{-}pos\text{-}term\ q\ (ss!i)$
by $auto$
with i **show** $?case$ **unfolding** $s\ ctxt\text{-}of\text{-}pos\text{-}term.simps\ to\text{-}pterm\text{-}ctxt.simps\ source\text{-}ctxt.simps$
using $source\text{-}to\text{-}pterm$ **by** $(metis\ source.simps(2)\ term.sel(4)\ to\text{-}pterm.simps(2))$

qed $simp$

lemma $to\text{-}pterm\text{-}ctxt\text{-}comp$: $to\text{-}pterm\text{-}ctxt\ (C \circ_c D) = to\text{-}pterm\text{-}ctxt\ C \circ_c to\text{-}pterm\text{-}ctxt\ D$
by$(induct\ C)$ $simp\text{-}all$

lemma $source\text{-}apply\text{-}subst$:
assumes $A \in wf\text{-}pterm\ R$
shows $source\ (A \cdot \sigma) = (source\ A) \cdot (source \circ \sigma)$
using $assms$ **proof**$(induct\ A)$
case $(3\ \alpha\ As)$
have $id{:}\forall\ x \in vars\text{-}term\ (lhs\ \alpha).\ ((\langle map\ (source \circ (\lambda t.\ t \cdot \sigma))\ As\rangle_\alpha)\ x = ((\langle map\ source\ As\rangle_\alpha \circ_s (source \circ \sigma))\ x$
proof−
have $vars{:}vars\text{-}term\ (lhs\ \alpha) = set\ (var\text{-}rule\ \alpha)$ **by** $simp$
{ fix i **assume** $i{:}i < length\ (var\text{-}rule\ \alpha)$
with 3 **have** $(\langle map\ (source \circ (\lambda t.\ t \cdot \sigma))\ As\rangle_\alpha)\ ((var\text{-}rule\ \alpha)!i) = source\ ((As!i) \cdot \sigma)$
by $(simp\ add{:}\ mk\text{-}subst\text{-}distinct)$
also have ... $= source\ (As!i) \cdot (source \circ \sigma)$
using $3\ i$ **by** $(metis\ nth\text{-}mem)$
also have ... $= (\langle map\ source\ As\rangle_\alpha \circ_s (source \circ \sigma))\ ((var\text{-}rule\ \alpha)!i)$
using $3\ i$ **unfolding** $subst\text{-}compose\text{-}def$ **by** $(simp\ add{:}\ mk\text{-}subst\text{-}distinct)$
finally have $(\langle map\ (source \circ (\lambda t.\ t \cdot \sigma))\ As\rangle_\alpha)\ ((var\text{-}rule\ \alpha)!i) = (\langle map\ source\ As\rangle_\alpha \circ_s (source \circ \sigma))\ ((var\text{-}rule\ \alpha)!i)$ **.**
} with $vars$ **show** $?thesis$ **by** $(metis\ in\text{-}set\text{-}idx)$
qed
have $source\ ((Prule\ \alpha\ As) \cdot \sigma) = (lhs\ \alpha) \cdot \langle map\ (source \circ (\lambda t.\ t \cdot \sigma))\ As\rangle_\alpha$
unfolding $eval\text{-}term.simps(2)$ **by** $simp$
also have ... $= (lhs\ \alpha) \cdot (\langle map\ source\ As\rangle_\alpha \circ_s (source \circ \sigma))$
using id **by** $(meson\ term\text{-}subst\text{-}eq)$
also have ... $= (source\ (Prule\ \alpha\ As)) \cdot (source \circ \sigma)$ **by** $simp$
finally show $?case$ **.**
qed $simp\text{-}all$

lemma $ctxt\text{-}of\text{-}pos\text{-}term\text{-}at\text{-}var\text{-}subst$:
assumes $linear\text{-}term\ t$
and $p \in poss\ t$ **and** $t|\text{-}p = Var\ x$
and $\forall\, y \in vars\text{-}term\ t.\ y \neq x \longrightarrow \tau\ y = \sigma\ y$
shows $ctxt\text{-}of\text{-}pos\text{-}term\ p\ (t \cdot \tau) = ctxt\text{-}of\text{-}pos\text{-}term\ p\ (t \cdot \sigma)$

using *assms* **proof**(*induct t arbitrary*:*p*)
case (*Fun f ts*)
from *Fun*(*3*,*4*) **obtain** *i* p' **where** *p*:$p = i\#p'$ **and** *i*:$i <$ *length ts* **and** p':$p' \in$ *poss* (*ts*!*i*)
by *auto*
with *Fun*(*4*) **have** *x*:*ts*!*i* |-p' = *Var x*
by *simp*
{**fix** *j* **assume** *j*:$j <$ *length ts* $j \neq i$
from *Fun*(*2*) **have** $x \notin$ *vars-term* (*ts*!*j*)
by (*metis i j* p' *subset-eq term.set-intros*(*3*) *var-in-linear-args vars-term-subt-at x*)
with *Fun*(*5*) *j* **have** *ts*!*j* $\cdot \tau$ = *ts*!*j* $\cdot \sigma$
by (*metis* (*no-types*, *lifting*) *nth-mem term.set-intros*(*4*) *term-subst-eq*)
then have (*map* ($\lambda t.\ t \cdot \tau$) *ts*)!*j* = (*map* ($\lambda t.\ t \cdot \sigma$) *ts*)!*j*
by (*simp add*: *j*)
}**note** *args*=*this*
from *args* **have** *args1*:*take i* (*map* ($\lambda t.\ t \cdot \tau$) *ts*) = *take i* (*map* ($\lambda t.\ t \cdot \sigma$) *ts*)
using *nth-take-lemma*[*of i* (*map* ($\lambda t.\ t \cdot \tau$) *ts*) (*map* ($\lambda t.\ t \cdot \sigma$) *ts*)] *i* **by** *simp*
from *args* **have** *args2*:*drop* (*Suc i*) (*map* ($\lambda t.\ t \cdot \tau$) *ts*) = *drop* (*Suc i*) (*map* ($\lambda t.\ t \cdot \sigma$) *ts*)
using *nth-drop-equal*[*of* (*map* ($\lambda t.\ t \cdot \tau$) *ts*) (*map* ($\lambda t.\ t \cdot \sigma$) *ts*) *Suc i*] *i* **by** *simp*
from *Fun*(*1*,*2*,*5*) *i* **have** *IH*:*ctxt-of-pos-term* p' ((*ts*!*i*) $\cdot \tau$) = *ctxt-of-pos-term* p' ((*ts*!*i*) $\cdot \sigma$)
by (*simp add*: p' *x*)
with *args1 args2* **show** *?case*
unfolding *p eval-term.simps ctxt-of-pos-term.simps* **by** (*simp add*: *i*)
qed *simp*

context *left-lin*
begin

lemma *source-ctxt-apply-subst*:
assumes $C \in$ *wf-pterm-ctxt R*
shows *source-ctxt* ($C \cdot_c \sigma$) = (*source-ctxt C*) $\cdot_c$ (*source* $\circ\ \sigma$)
using *assms* **proof**(*induct C*)
case (*2 ss1 ss2 C f*)
then show *?case*
unfolding *source-ctxt.simps actxt.simps 2* **using** *source-apply-subst* **by** *auto*
next
case (*3* α *ss1 ss2 C*)
let *?p*=(*var-poss-list* (*lhs* α) ! *length ss1*)
let *?x*=(*vars-term-list* (*lhs* α) ! *length ss1*)
have *var-at-p*:(*lhs* α)|-*?p* = *Var ?x*
by (*metis 3.hyps*(*2*) *add-lessD1 length-remdups-leq length-rev less-add-one o-apply order-less-le-trans vars-term-list-var-poss-list*)
from *3*(*2*) **have** *pos1*:*?p* $\in$ *poss* (*lhs* α)
by (*metis add-lessD1 comp-apply length-remdups-leq length-rev length-var-poss-list less-add-one nth-mem order-less-le-trans var-poss-imp-poss var-poss-list-sound*)

then have *pos*:*?p* $\in$ *poss* (*lhs* $\alpha \cdot \langle$*map source* (*ss1* @ *Var* (*vars-term-list* (*lhs* α) ! *length ss1*) # *ss2*)$\rangle_\alpha$)
using *poss-imp-subst-poss* **by** *blast*
have *lin*:*linear-term* (*lhs* α)
using *3*(*1*) *left-lin* **using** *left-linear-trs-def* **by** *fastforce*
{**fix** *y* **assume** *y* $\in$ *vars-term* (*lhs* α) **and** *x*:*y* $\neq$ *?x*
then obtain *i* **where** *i*:*i* $<$ *length* (*var-rule* α) *var-rule* α ! *i* = *y*
by (*metis in-set-idx lin linear-term-var-vars-term-list set-vars-term-list*)
with *x* **consider** *i* $<$ *length ss1* | *i* $>$ *length ss1* $\wedge$ *i* $<$ *length* (*var-rule* α)
using *lin linear-term-var-vars-term-list nat-neq-iff* **by** *fastforce*
then have ($\langle$*map source* (*map* ($\lambda t.\ t \cdot \sigma$) *ss1* @ *Var ?x* # *map* ($\lambda t.\ t \cdot \sigma$) *ss2*)$\rangle_\alpha$) *y* = (($\langle$*map source* (*ss1* @ *Var ?x* # *ss2*)$\rangle_\alpha$) *y*) $\cdot$ (*source* $\circ$ σ)
proof(*cases*)
case *1*
with *i* **have** ($\langle$*map source* (*map* ($\lambda t.\ t \cdot \sigma$) *ss1* @ *Var ?x* # *map* ($\lambda t.\ t \cdot \sigma$) *ss2*)$\rangle_\alpha$) *y* = *source* ((*ss1*!*i*) $\cdot$ σ)
by (*smt* (*z3*) *3.hyps*(*2*) *One-nat-def add.right-neutral add-Suc-right append-Cons-nth-left comp-apply distinct-remdups distinct-rev length-append length-map length-nth-simps*(*2*) *map-nth-eq-conv mk-subst-same*)
moreover from *i 1* **have** ($\langle$*map source* (*ss1* @ *Var ?x* # *ss2*)$\rangle_\alpha$) *y* = *source* (*ss1*!*i*)
by (*smt* (*verit*, *ccfv-threshold*) *3.hyps*(*2*) *One-nat-def ab-semigroup-add-class.add-ac*(*1*) *append-Cons-nth-left comp-apply distinct-remdups distinct-rev length-append length-map list.size*(*4*) *map-nth-eq-conv mk-subst-distinct*)
moreover have *ss1*!*i* $\in$ *wf-pterm R*
using *3*(*3*) *1* **by** (*meson UnCI nth-mem*)
ultimately show *?thesis*
using *source-apply-subst* **by** *auto*
next
case *2*
let *?i*=*i* $-$ ((*length ss1*)+*1*)
have *i'*:*?i* $<$ *length ss2*
using *3*(*2*) *2* **by** (*simp add*: *less-diff-conv2*)
have *i1*:(*map source* (*map* ($\lambda t.\ t \cdot \sigma$) *ss1* @ *Var ?x* # *map* ($\lambda t.\ t \cdot \sigma$) *ss2*))!*i* = *source* ((*ss2*!*?i*) $\cdot$ σ) **proof**$-$
have *i''*:*i* = *length* (*map source* (*map* ($\lambda t.\ t \cdot \sigma$) *ss1* @ [*Var ?x*])) + *?i*
unfolding *length-append length-map* **using** *2* **by** *force*
show *?thesis* **unfolding** *map-append list.map*
using *i' i''* *nth-append-length-plus*[*of* (*map source* (*map* ($\lambda t.\ t \cdot \sigma$) *ss1* @ [*Var* (*vars-term-list* (*lhs* α) ! *length ss1*)])) *map source* (*map* ($\lambda t.\ t \cdot \sigma$) *ss2*)]
by (*smt* (*verit*, *del-insts*) *Cons-eq-appendI append-Nil append-assoc length-map list.simps*(*9*) *map-append nth-map*)
qed
have *i2*:*map source* (*ss1* @ *Var ?x* # *ss2*) ! *i* = *source* (*ss2*!*?i*) **proof**$-$
have *i''*:*i* = *length* (*map source* (*ss1* @ [*Var ?x*])) + *?i*
unfolding *length-append length-map* **using** *2* **by** *force*
show *?thesis* **unfolding** *map-append list.map*
using *i' i''* *nth-append-length-plus*[*of* (*map source* (*ss1* @ [*Var* (*vars-term-list* (*lhs* α) ! *length ss1*)])) *map source ss2*]

by (*smt* (*verit*, *del-insts*) *append.left-neutral append-Cons append-assoc length-map list.simps*(*9*) *map-append nth-map*)
qed
from *i1 2* **have** ($\langle$*map source* (*map* ($\lambda t.\ t \cdot \sigma$) *ss1* @ *Var ?x* # *map* ($\lambda t.\ t \cdot \sigma$) *ss2*)$\rangle_\alpha$) *y* = *source* ((*ss2*!*?i*) $\cdot$ σ)
by (*smt* (*verit*, *ccfv-threshold*) *3.hyps*(*2*) *One-nat-def ab-semigroup-add-class.add-ac*(*1*) *comp-def distinct-remdups distinct-rev i*(*2*) *length-append length-map list.size*(*4*) *mk-subst-distinct*)
moreover from *i2 2* **have** ($\langle$*map source* (*ss1* @ *Var ?x* # *ss2*)$\rangle_\alpha$) *y* = *source* (*ss2*!*?i*)
by (*metis* (*no-types*, *opaque-lifting*) *3.hyps*(*2*) *One-nat-def add.right-neutral add-Suc-right comp-apply distinct-remdups distinct-rev i*(*2*) *length-append length-map length-nth-simps*(*2*) *mk-subst-distinct*)
moreover have *ss2*!*?i* $\in$ *wf-pterm R*
using *3*(*3*) *2* ‹*?i* < *length ss2*› **by** (*metis UnCI nth-mem*)
ultimately show *?thesis*
using *source-apply-subst* **by** *auto*
qed
}
then have *ctxt-of-pos-term ?p* (*lhs* α $\cdot$ $\langle$*map source* (*map* ($\lambda t.\ t \cdot \sigma$) *ss1* @ *Var ?x* # *map* ($\lambda t.\ t \cdot \sigma$) *ss2*)$\rangle_\alpha$) =
ctxt-of-pos-term ?p (*lhs* α $\cdot$ $\langle$*map source* (*ss1* @ *Var ?x* # *ss2*)$\rangle_\alpha$ $\cdot$ (*source* $\circ$ σ))
using *ctxt-of-pos-term-at-var-subst*[*OF lin pos1 var-at-p*] **unfolding** *subst-subst* **by** (*smt* (*verit*) *subst-compose*)
then show *?case* **unfolding** *source-ctxt.simps actxt.simps Let-def 3 subst-compose-ctxt-compose-distrib length-map ctxt-of-pos-term-subst*[*OF pos*, *symmetric*]
by *presburger*
qed *simp*

Needs left-linearity to avoid multihole contexts.

lemma *source-ctxt-apply-term*:
assumes *C* $\in$ *wf-pterm-ctxt R*
shows *source* (*C*$\langle$*A*$\rangle$) = (*source-ctxt C*)$\langle$*source A*$\rangle$
using *assms* **proof**(*induct C*)
case (*3* α *ss1 ss2 C*)
from *3*(*1*) *left-lin* **have** *lin*:*linear-term* (*lhs* α)
using *left-linear-trs-def* **by** *fastforce*
from *3*(*2*) **have** *len*:*length ss1* < *length* (*vars-term-list* (*lhs* α))
by (*metis add-lessD1 less-add-one lin linear-term-var-vars-term-list*)
have (*source-ctxt* (*Crule* α *ss1 C ss2*))$\langle$*source A*$\rangle$ =
lhs α $\cdot$ $\langle$(*map source ss1*) @ (*source-ctxt C*)$\langle$*source A*$\rangle$ # (*map source ss2*)$\rangle_\alpha$
unfolding *source-ctxt.simps Let-def intp-actxt.simps source.simps ctxt-ctxt-compose*
using *ctxt-apply-term-subst*[*OF lin len*] *lhs-subst-upd*
by (*smt* (*verit*) *len length-map lin linear-term-var-vars-term-list list.simps*(*9*) *map-append*)
with *3*(*5*) **show** *?case* **by** *simp*
qed *simp-all*
end

lemma *rewrite-tgt*:
  assumes *rstep*:$(t,v) \in (rstep\ R)^*$
  shows $(target\ (C\ \langle(to\text{-}pterm\ t) \cdot \sigma\rangle),\ target\ (C\ \langle(to\text{-}pterm\ v) \cdot \sigma\rangle)) \in (rstep\ R)^*$
proof(*induct C*)
  case *Hole*
  then show *?case*
   by (*simp add*: *local.rstep rsteps-closed-subst target-empty-apply-subst to-pterm-empty*)

next
  case (*Cfun f ss1 C ss2*)
  then show *?case* **by** (*simp add*: *ctxt-closed-one ctxt-closed-rsteps*)
next
  case ($Crule\ \alpha\ ss1\ C\ ss2$)
  {**fix** x **assume** $x \in vars\text{-}term\ (rhs\ \alpha)$
    from *Crule* **have** $((\langle map\ target\ (ss1\ @\ C\langle to\text{-}pterm\ t \cdot \sigma\rangle\ \#\ ss2)\rangle_\alpha)\ x,\ (\langle map\ target\ (ss1\ @\ C\langle to\text{-}pterm\ v \cdot \sigma\rangle\ \#\ ss2)\rangle_\alpha)\ x) \in (rstep\ R)^*$
    proof($cases\ x \in vars\text{-}term\ (lhs\ \alpha)$)
      case *True*
     then obtain i **where** $i < length\ (vars\text{-}distinct\ (lhs\ \alpha))\ x = vars\text{-}distinct\ (lhs\ \alpha)!i$
        by (*metis in-set-idx set-vars-term-list vars-term-list-vars-distinct*)
      then show *?thesis* **using** *Crule*
      by (*smt* (*z3*) *append-Cons-nth-not-middle length-append length-map length-nth-simps*(*2*) *lhs-subst-not-var-i lhs-subst-var-i map-nth-eq-conv nth-append-length rtrancl.simps*)

    next
      case *False*
      then show *?thesis*
        by (*simp add*: *mk-subst-not-mem*)
    qed
  }
  then show *?case* **by** (*simp add*: *subst-rsteps-imp-rsteps*)
qed

2.8 Additional Results

lemma *length-args-well-Prule*:
  assumes $Prule\ \alpha\ As \in wf\text{-}pterm\ R\ Prule\ \alpha\ Bs \in wf\text{-}pterm\ S$
  shows $length\ As = length\ Bs$
proof−
  from *assms*(*1*) **have** $length\ As = length\ (var\text{-}rule\ \alpha)$ **using** *wf-pterm.simps* **by** *fastforce*
 moreover from *assms*(*2*) **have** $length\ Bs = length\ (var\text{-}rule\ \alpha)$ **using** *wf-pterm.simps* **by** *fastforce*
  ultimately show *?thesis* **by** *simp*
qed

lemma *co-initial-Var*:

assumes *co-initial* (*Var* x) B
shows $B = Var\ x \vee (\exists\, \alpha\ b'\ y.\ B = Prule\ \alpha\ b' \wedge lhs\ \alpha = Var\ y)$
proof−
{**assume** $B \neq Var\ x$
with *assms* **obtain** α b' **where** $B = Prule\ \alpha\ b'$
by (*metis is-empty-step.elims*(*3*) *source.elims source-empty-step term.distinct*(*1*))
moreover with *assms* **have** $\exists\ y.\ lhs\ \alpha = Var\ y$
by (*metis source.simps*(*1*) *source.simps*(*3*) *subst-apply-eq-Var*)
ultimately have $(\exists\, \alpha\ b'\ y.\ B = Prule\ \alpha\ b' \wedge lhs\ \alpha = Var\ y)$
by *blast*
}
then show *?thesis*
by *blast*
qed

lemma *source-poss*:
assumes p:$p \in poss\ (source\ (Pfun\ f\ As))$ **and** iq:$i\#q \in poss\ (Pfun\ f\ As)$
and *ctxt*:*source-ctxt* (*ctxt-of-pos-term* ($i\#q$) (*Pfun f As*)) = *ctxt-of-pos-term* p (*source* (*Pfun f As*))
shows $\exists\, p'.\ p = i\#p' \wedge p' \in poss\ (source\ (As!i))$
proof−
obtain p' **where** *hole-pos* (*source-ctxt* (*ctxt-of-pos-term* ($i\#q$) (*Pfun f As*))) = $i\#p'$
p' = *hole-pos* (*source-ctxt* (*ctxt-of-pos-term* q (*As* ! i)))
unfolding *ctxt-of-pos-term.simps source-ctxt.simps take-map drop-map* **using** *iq* **by** *auto*
with *ctxt* **have** $p = i\#p'$
by (*metis hole-pos-ctxt-of-pos-term p*)
with p **show** *?thesis*
by *auto*
qed

lemma *simple-pterm-match*:
assumes *source* $A = t \cdot \sigma$
and *linear-term* t
and $A \cdot \tau 1 = to\text{-}pterm\ t \cdot \tau 2$
shows *matches* A (*to-pterm* t)
using *assms* **proof**(*induct t arbitrary*: A)
case (*Var* x)
then show *?case*
using *matches-iff* **by** *force*
next
case (*Fun f ts*)
from *Fun*(*2*,*4*) **show** *?case* **proof**(*cases A*)
case (*Pfun g As*)
with *Fun*(*2*) **have** f:$f = g$ **by** *simp*
from *Fun*(*2*) **have** l:*length ts* = *length As*
unfolding *Pfun source.simps f eval-term.simps* **by** (*simp add*: *map-equality-iff*)

{**fix** *i* **assume** *i*:*i* < *length ts*
with *Fun*(*2*) **have** *source* (*As* ! *i*) = *ts* ! *i* · σ
unfolding *Pfun source.simps f eval-term.simps* **by** (*simp add*: *map-equality-iff*)
moreover from *i Fun*(*4*) **have** *As* ! *i* · τ*1* = *to-pterm* (*ts* ! *i*) · τ*2*
unfolding *Pfun f to-pterm.simps eval-term.simps* **using** *l map-nth-conv* **by** *fastforce*
ultimately have *matches* (*As*!*i*) (*to-pterm* (*ts*!*i*))
using *Fun*(*1*)[*of ts*!*i As*!*i*] *l i Fun*(*3*) **by** *force*
then have ∃σ. *As*!*i* = (*to-pterm* (*ts*!*i*)) · σ
by (*metis matches-iff*)
}**note** *IH*=*this*
from *Fun*(*3*) **have** *lin*:*linear-term* (*to-pterm* (*Fun f ts*))
using *to-pterm-linear* **by** *blast*
from *linear-term-obtain-subst*[*OF lin*[*unfolded to-pterm.simps*]] **show** *?thesis*
unfolding *Pfun f* **by** (*smt* (*verit*, *del-insts*) *IH l length-map matches-iff nth-map to-pterm.simps*(*2*))
qed *simp-all*
qed

2.9 Proof Terms Represent Multi-Steps

context *var-rhs-subset-lhs*
begin
lemma *mstep-to-pterm*:
assumes (*s*, *t*) ∈ *mstep R*
shows ∃*A*. *A* ∈ *wf-pterm R* ∧ *source A* = *s* ∧ *target A* = *t*
using *assms*(*1*) **proof**(*induct*)
case (*Var x*)
then show *?case*
by (*meson source.simps*(*1*) *target.simps*(*1*) *wf-pterm.intros*(*1*))
next
case (*args f n ss ts*)
then have ∀ *i*∈*set* [*0*..<*n*]. ∃ *a*. *a* ∈ *wf-pterm R* ∧ *source a* = *ss* ! *i* ∧ *target a* = *ts* ! *i*
by *simp*
then obtain *As* **where** *as*:*length As* = *n* ∧ (∀ *i* < *n*. (*As*!*i*) ∈ *wf-pterm R* ∧ *source* (*As*!*i*) = *ss* ! *i* ∧ *target* (*As*!*i*) = *ts* ! *i*)
using *obtain-list-with-property*[**where** *P*=λ*a i*. *a* ∈ *wf-pterm R* ∧ *source a* = *ss*!*i* ∧ *target a* = *ts*!*i* **and** *xs*=[*0*..<*n*]]
by (*metis add.left-neutral diff-zero length-upt nth-upt set-upt*)
with *args*(*1*) **have** *source* (*Pfun f As*) = *Fun f ss*
unfolding *source.simps* **by** (*simp add*: *map-nth-eq-conv*)
moreover from *as args*(*2*) **have** *target* (*Pfun f As*) = *Fun f ts*
unfolding *target.simps* **by** (*simp add*: *map-nth-eq-conv*)
ultimately show *?case*
using *as* **by** (*metis in-set-idx wf-pterm.intros*(*2*))
next
case (*rule l r* σ τ)
let *?*α =(*l* → *r*)

have *set* (*vars-distinct l*) = *vars-term l*
by *simp*
with *rule*(*2*) **obtain** *As* **where** *as*:*length As* = *length* (*vars-distinct l*) ∧
(∀ *i* < *length* (*vars-distinct l*). (*As*!*i*) ∈ *wf-pterm R* ∧
source (*As*!*i*) = σ ((*vars-distinct l*) ! *i*) ∧ *target* (*As*!*i*) = τ ((*vars-distinct l*) ! *i*))
using *obtain-list-with-property*[**where** *P*=λ*a x*. *a* ∈ *wf-pterm R* ∧ *source a* = σ *x* ∧ *target a* = τ *x*] **by** *blast*
with *rule*(*1*) **have** *well*:*Prule ?α As* ∈ *wf-pterm R*
by (*metis in-set-idx prule.sel*(*1*) *prule.sel*(*2*) *wf-pterm.simps*)
from *as* **have** ∀ *x* ∈ *vars-term l*. (⟨*map source As*⟩$_{?}$α) *x* = σ *x*
by (*smt* (*z3*) *apply-lhs-subst-var-rule in-set-idx length-map map-nth-conv prule.sel*(*1*) *set-vars-term-list vars-term-list-vars-distinct*)
then have *s*:*source* (*Prule ?α As*) = *l* · σ
by (*simp add*: *term-subst-eq-conv*)
from *as varcond* **have** ∀ *x* ∈ *vars-term r*. (⟨*map target As*⟩$_{?}$α) *x* = τ *x*
by (*smt* (*verit*, *best*) *apply-lhs-subst-var-rule fst-conv in-set-conv-nth length-map nth-map prule.sel*(*1*)
rule.hyps(*1*) *set-vars-term-list snd-conv split-beta subsetD vars-term-list-vars-distinct*)

then have *target* (*Prule ?α As*) = *r* · τ
by (*simp add*: *term-subst-eq-conv*)
with *well s* **show** *?case*
by *blast*
qed
end

lemma *pterm-to-mstep*:
assumes *A* ∈ *wf-pterm R*
shows (*source A*, *target A*) ∈ *mstep R*
using *assms* **proof**(*induct*)
case (*2 As f*)
then show *?case*
by (*simp add*: *mstep.args*)
next
case (*3* α *As*)
then have ∀ *x*∈*vars-term* (*lhs* α). (((⟨*map source As*⟩$_{α}$) *x*, (⟨*map target As*⟩$_{α}$) *x*) ∈ *mstep R*
by (*smt* (*verit*, *best*) *apply-lhs-subst-var-rule comp-def in-set-idx length-map map-nth-conv nth-mem set-remdups set-rev set-vars-term-list*)
with *3*(*1*) **show** *?case*
by (*simp add*: *mstep.rule*)
qed *simp*

lemma *co-init-prule*:
assumes *co-initial* (*Prule* α *As*) (*Prule* α *Bs*)
and *Prule* α *As* ∈ *wf-pterm R* **and** *Prule* α *Bs* ∈ *wf-pterm R*
shows ∀ *i*<*length As*. *co-initial* (*As*!*i*) (*Bs*!*i*)
proof−

from *assms* **have** $l1$:*length As* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
from *assms* **have** $l2$:*length Bs* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
{**fix** i **assume** i:i < *length As* **and** *co*:$\neg$ (*co-initial* (*As*!i) (*Bs*!i))
then have $(\langle map\ source\ As\rangle_\alpha)$ ((*var-rule* α)!i) $\neq$ $(\langle map\ source\ Bs\rangle_\alpha)$ ((*var-rule* α)!i)
by (*metis l1 l2 length-map lhs-subst-var-i nth-map*)
with *assms*(*1*) **have** *False* **unfolding** *source.simps*
by (*smt* (*z3*) *comp-apply i l1 nth-mem set-remdups set-rev set-vars-term-list term-subst-eq-rev*)
} **then show** *?thesis*
by *blast*
qed

3 Operations on Proof Terms

The operations residual, deletion, and join on proof terms all fulfill $A \star (source\ A) = A$ which implies several useful results.

locale *op-proof-term* = *left-lin-no-var-lhs* +
fixes f :: (('*a*, '*b*) *prule* + '*a*, '*b*) *Term.term* $\Rightarrow$ (('*a*, '*b*) *prule* + '*a*, '*b*) *Term.term* $\Rightarrow$ (('*a*, '*b*) *prule* + '*a*, '*b*) *Term.term option*
assumes *f-src*: $A \in$ *wf-pterm R* $\Longrightarrow$ *f A* (*to-pterm* (*source A*)) = *Some A*
and *f-pfun*:*f* (*Pfun g As*)(*Pfun g Bs*) = (*if length As* = *length Bs then*
(*case those* (*map2 f As Bs*) *of*
Some xs $\Rightarrow$ *Some* (*Pfun g xs*)
| *None* $\Rightarrow$ *None*) *else None*)
and *f-prule*:*f* (*Prule* α *As*) (*Pfun g Bs*) = (*case match* (*Pfun g Bs*) (*to-pterm* (*lhs* α)) *of*
None $\Rightarrow$ *None*
| *Some* σ $\Rightarrow$
(*case those* (*map2 f As* (*map* σ (*var-rule* α))) *of*
Some xs $\Rightarrow$ *Some* (*Prule* α *xs*)
| *None* $\Rightarrow$ *None*))
begin

notation
f ('($\star$')) **and**
f ((- $\star$ -) [*51*, *51*] *50*)

lemma *apply-f-ctxt*:
assumes $C \in$ *wf-pterm-ctxt R*
and $A \star B$ = *Some D*
shows $C\langle A\rangle \star$ (*to-pterm-ctxt* (*source-ctxt C*))$\langle B\rangle$ = *Some* ($C\langle D\rangle$)
using *assms* **proof**(*induct C rule*:*pterm-ctxt-induct*)
case (*Cfun f ss1 C ss2*)
have *l*:*length* ((*map* (*to-pterm* $\circ$ *source*) *ss1*) @ (*to-pterm-ctxt* (*source-ctxt C*))$\langle A\rangle$ # (*map* (*to-pterm* $\circ$ *source*) *ss2*))

$= length\ (ss1\ @\ C\langle B\rangle\ \#\ ss2)$ **by** *auto*
from *Cfun*(*2*) **have** $well1{:}\forall\, i < length\ ss1.\ (ss1!i) \in wf\text{-}pterm\ R$ **by** *auto*
from *Cfun*(*2*) **have** $well2{:}\forall\, i < length\ ss2.\ (ss2!i) \in wf\text{-}pterm\ R$ **by** *auto*
from *Cfun* **have** $fC{:}C\langle A\rangle \star (to\text{-}pterm\text{-}ctxt\ (source\text{-}ctxt\ C))\langle B\rangle = Some\ (C\langle D\rangle)$

by *auto*
from *well1* **have** $f1{:}\forall\, i < length\ ss1.\ ((map2\ (\star)\ ss1\ (map\ (to\text{-}pterm \circ source)\ ss1))!i = Some\ (ss1!i))$
using *f-src to-pterm-empty* **by** *fastforce*
from *well2* **have** $f2{:}\forall\, i < length\ ss2.\ ((map2\ (\star)\ ss2\ (map\ (to\text{-}pterm \circ source)\ ss2))!i = Some\ (ss2!i))$
using *f-src to-pterm-empty* **by** *fastforce*
{fix *i* **assume** $i{:}i < (length\ ss1) + (length\ ss2) + 1$
have $(map2\ (\star)\ (ss1\ @\ (C\langle A\rangle\ \#\ ss2))$
$(map\ (to\text{-}pterm \circ source)\ ss1\ @\ ((to\text{-}pterm\text{-}ctxt\ (source\text{-}ctxt\ C))\langle B\rangle\ \#\ map\ (to\text{-}pterm \circ source)\ ss2)))!i$
$= Some\ ((ss1\ @\ C\langle D\rangle\ \#\ ss2)!i)$
proof−
consider $i < length\ ss1 \mid i = length\ ss1 \mid i > length\ ss1$
using *nat-neq-iff* **by** *blast*
then show *?thesis* **proof**(*cases*)
case *1*
then show *?thesis* **using** *f1*
by (*simp add*: *append-Cons-nth-left*)
next
case *2*
then show *?thesis* **using** *fC*
by (*simp add*: *append-Cons-nth-middle*)
next
case *3*
with *i* **have** $l{:}(map\ (to\text{-}pterm \circ source)\ ss1\ @\ (to\text{-}pterm\text{-}ctxt\ (source\text{-}ctxt\ C))\langle B\rangle\ \#\ map\ (to\text{-}pterm \circ source)\ ss2)!i$
$= (map\ (to\text{-}pterm \circ source)\ ss2)!(i-(length\ ss1 + 1))$
by (*metis add.commute length-map less-SucI not-less-eq nth-append-Cons plus-1-eq-Suc*)
from *3 i* **have** $r{:}(ss1\ @\ (C\langle to\text{-}pterm\ (source\ B)\rangle\ \#\ ss2))!i = ss2!(i-(length\ ss1 + 1))$
by (*metis add.commute less-SucI not-less-eq nth-append-Cons plus-1-eq-Suc*)

from *l r 3* **show** *?thesis* **using** *f2*
by (*smt One-nat-def add.right-neutral add-Suc add-Suc-right add-diff-inverse-nat add-less-cancel-left append-Cons-nth-right i length-append length-map length-zip list.size*(*4*) *min-less-iff-conj not-less-eq nth-map nth-zip*)
qed
qed
}
with *l* **have** $those\ ((map2\ (\star)\ (ss1\ @\ (C\langle A\rangle\ \#\ ss2))$
$(map\ (to\text{-}pterm \circ source)\ ss1\ @\ ((to\text{-}pterm\text{-}ctxt\ (source\text{-}ctxt\ C))\langle B\rangle\ \#\ map\ (to\text{-}pterm \circ source)\ ss2)))$

$= Some\ (ss1\ @\ C\langle D\rangle\ \#\ ss2)$ **by** *(simp add: those-some)*
with *l* **show** *?case* **using** *f-pfun* **by** *simp*
next
case *(Crule α ss1 C ss2)*
from *Crule(2)* **have** *alpha:to-rule* $\alpha \in R$
using *wf-pterm-ctxt.cases* **by** *auto*
then have *linear:linear-term (lhs α)*
using *left-lin left-linear-trs-def* **by** *fastforce*
then have *linear':linear-term (to-pterm (lhs α))*
using *to-pterm-linear* **by** *blast*
have *l1:length ((map (to-pterm $\circ$ source) ss1) @ (to-pterm-ctxt (source-ctxt C))$\langle A\rangle$ # (map (to-pterm $\circ$ source) ss2))*
$= length\ (ss1\ @\ C\langle B\rangle\ \#\ ss2)$ **by** *auto*
from *Crule(2)* **have** *l2:length (ss1 @ C$\langle B\rangle$ # ss2) = length (var-rule α)*
using *wf-pterm-ctxt.simps* **by** *fastforce*
from *Crule(2)* **have** *well1:*$\forall\, i < length\ ss1.\ (ss1!i) \in wf\text{-}pterm\ R$ **by** *auto*
from *Crule(2)* **have** *well2:*$\forall\, i < length\ ss2.\ (ss2!i) \in wf\text{-}pterm\ R$ **by** *auto*
from *Crule* **have** *fC:C$\langle A\rangle$ $\star$ (to-pterm-ctxt (source-ctxt C))$\langle B\rangle$ = Some (C$\langle D\rangle$)*

by *auto*
from *well1* **have** *f1:*$\forall\, i < length\ ss1.$ *((map2 ($\star$) ss1 (map (to-pterm $\circ$ source) ss1))!i = Some (ss1!i))*
using *f-src to-pterm-empty* **by** *fastforce*
from *well2* **have** *f2:*$\forall\, i < length\ ss2.$ *((map2 ($\star$) ss2 (map (to-pterm $\circ$ source) ss2))!i = Some (ss2!i))*
using *f-src to-pterm-empty* **by** *fastforce*
{**fix** *i* **assume** *i:i < (length ss1) + (length ss2) +1*
have *(map2 ($\star$) (ss1 @ (C$\langle A\rangle$ # ss2)) (map (to-pterm $\circ$ source) ss1 @ ((to-pterm-ctxt (source-ctxt C))$\langle B\rangle$ #*
(map (to-pterm $\circ$ source) ss2))))!i = Some ((ss1 @ C$\langle D\rangle$ # ss2)!i)
proof−
consider *i < length ss1 | i = length ss1 | i > length ss1*
using *nat-neq-iff* **by** *blast*
then show *?thesis* **proof**(*cases*)
case *1*
then show *?thesis* **using** *f1*
by *(simp add: append-Cons-nth-left)*
next
case *2*
then show *?thesis* **using** *fC*
by *(simp add: append-Cons-nth-middle)*
next
case *3*
with *i* **have** *l:(map (to-pterm $\circ$ source) ss1 @ (to-pterm-ctxt (source-ctxt C))$\langle B\rangle$ # map (to-pterm $\circ$ source) ss2)!i*
= (map (to-pterm $\circ$ source) ss2)!(i−(length ss1 + 1))
by *(metis add.commute length-map less-SucI not-less-eq nth-append-Cons plus-1-eq-Suc)*
from *3 i* **have** *r:(ss1 @ (C$\langle$to-pterm (source B)$\rangle$ # ss2))!i = ss2!(i−(length*

ss1 + *1*))
  by (*metis add.commute less-SucI not-less-eq nth-append-Cons plus-1-eq-Suc*)

  from *l r 3* **show** *?thesis* **using** *f2*
  by (*smt One-nat-def add.right-neutral add-Suc add-Suc-right add-diff-inverse-nat add-less-cancel-left append-Cons-nth-right i length-append length-map length-zip list.size*(*4*) *min-less-iff-conj not-less-eq nth-map nth-zip*)
  qed
  qed
  }
  with *l1* **have** *IH*:*those* (*map2* ($\star$) (*ss1* @ ($C\langle A\rangle$ # *ss2*)) (*map* (*to-pterm* $\circ$ *source*) *ss1* @ ((*to-pterm-ctxt* (*source-ctxt* *C*))$\langle B\rangle$ #
  (*map* (*to-pterm* $\circ$ *source*) *ss2*)))) = *Some* (*ss1* @ $C\langle D\rangle$ # *ss2*) **by** (*simp add*: *those-some*)
  let *?p* = (*var-poss-list* (*lhs* α) ! *length ss1*)
  let *?x* = *vars-term-list* (*lhs* α) ! *length ss1*
  let *?σ* = $\langle$*map source* (*ss1* @ *Var* (*vars-term-list* (*lhs* α) ! *length ss1*) # *ss2*)$\rangle_\alpha$
  from *l2 linear* **have** *l3*:*length ss1* < *length* (*var-poss-list* (*lhs* α))
  by (*metis* (*no-types*, *lifting*) *add-Suc-right append-Cons-nth-left le-imp-less-Suc length-append length-var-poss-list linear-term-var-vars-term-list linorder-neqE-nat list.size*(*3*) *list.size*(*4*) *not-add-less1 nth-equalityI self-append-conv zero-order*(*1*))
  then have *?p* $\in$ *poss* (*lhs* α)
  using *nth-mem var-poss-imp-poss var-poss-list-sound* **by** *blast*
  then have *ctxt*:(*to-pterm-ctxt* (*source-ctxt* (*Crule* α *ss1 C ss2*)))$\langle B\rangle$ =
  (*ctxt-of-pos-term* *?p* (*to-pterm* (*lhs* α) $\cdot$ (*to-pterm* $\circ$ *?σ*)))$\langle$(*to-pterm-ctxt* (*source-ctxt* *C*))$\langle B\rangle\rangle$
  unfolding *source-ctxt.simps intp-actxt.simps Let-def ctxt-ctxt-compose to-pterm-ctxt-comp*

  using *to-pterm-ctxt-at-pos*[**where** *?p*=*?p* **and** *?s*=*lhs* α $\cdot$ *?σ*] **by** (*simp add*: *to-pterm-subst*)
  from *l3* **have** *l4*:*length ss1* < *length* (*vars-term-list* (*to-pterm* (*lhs* α)))
  by (*metis length-var-poss-list vars-to-pterm*)
  have (*to-pterm-ctxt* (*source-ctxt* (*Crule* α *ss1 C ss2*)))$\langle B\rangle$ =
  to-pterm (*lhs* α) $\cdot$ ((*to-pterm* $\circ$ *?σ*)(*?x* := (*to-pterm-ctxt* (*source-ctxt* *C*))$\langle B\rangle$))
  unfolding *ctxt* **using** *ctxt-apply-term-subst*[**where** *?p*=*?p* **and** *?t*=*to-pterm* (*lhs* α) **and** *?i*=*length ss1* **and** *?s*=(*to-pterm-ctxt* (*source-ctxt* *C*))$\langle B\rangle$ **and** *?σ*=(*to-pterm* $\circ$ *?σ*)]
  linear′ l4 var-poss-list-to-pterm vars-to-pterm **by** *metis*
  then obtain τ **where** τ:*match* (*to-pterm-ctxt* (*source-ctxt* (*Crule* α *ss1 C ss2*)))$\langle B\rangle$ (*to-pterm* (*lhs* α)) = *Some* τ
  unfolding *ctxt* **using** *ctxt-apply-term-subst linear′ match-complete′ option.distinct*(*1*) **by** *force*
  have *varr*:(*var-rule* α) = *vars-term-list* (*to-pterm* (*lhs* α))
  using *linear linear-term-var-vars-term-list* **unfolding** *vars-to-pterm* **by** *force*
  have (*map* (*to-pterm* $\circ$ *?σ*) (*vars-term-list* (*to-pterm* (*lhs* α)))) = *map* (*to-pterm* $\circ$ *source*) (*ss1* @ *Var* (*vars-term-list* (*lhs* α) ! *length ss1*) # *ss2*)
  using *apply-lhs-subst-var-rule l2* **unfolding** *varr*[*symmetric*] **by** *force*
  then have (*map* (*to-pterm* $\circ$ *?σ*) (*vars-term-list* (*to-pterm* (*lhs* α))))[*length ss1*

:= (*to-pterm-ctxt* (*source-ctxt* C))$\langle B\rangle$] =
(*map* (*to-pterm* $\circ$ *source*) *ss1* @ (*to-pterm-ctxt* (*source-ctxt* C))$\langle B\rangle$ #
(*map* (*to-pterm* $\circ$ *source*) *ss2*))
by (*metis* (*no-types*, *lifting*) *length-map list.simps*(*9*) *list-update-length map-append*)

with τ **have** *map-tau*:*map* τ (*var-rule* α) = (*map* (*to-pterm* $\circ$ *source*) *ss1* @
(*to-pterm-ctxt* (*source-ctxt* C))$\langle B\rangle$ #
(*map* (*to-pterm* $\circ$ *source*) *ss2*))
using *match-lhs-context*[**where** *?t*=*to-pterm* (*lhs* α) **and** *?*τ=τ **and** *?*σ=(*to-pterm*
$\circ$ *?*σ)]
l4 var-poss-list-to-pterm linear′ ctxt varr **by** *metis*
from *alpha no-var-lhs* **obtain** *f ts* **where** *f*:*lhs* α = *Fun f ts*
by *blast*
have [] $\notin$ *var-poss* (*lhs* α)
unfolding *f var-poss.simps* **by** *force*
then obtain *i q* **where** *iq*:*?p* = *i* # *q* **using** *l3*
by (*metis in-set-conv-nth subt-at.elims var-poss-list-sound*)
then obtain *ts′* **where** *root-not-rule*:(*to-pterm-ctxt* (*source-ctxt* (*Crule* α *ss1* C
ss2)))$\langle B\rangle$ = *Pfun f ts′*
unfolding *ctxt iq* **unfolding** *f* **by** *simp*
then show *?case*
using τ *f-prule map-tau IH* **by** *force*
qed *simp*

end

end
theory *Residual-Join-Deletion*

imports
Proof-Terms
Linear-Matching
begin

3.1 Residuals

Auxiliary lemma in preparation of termination simp rule.

lemma *match-vars-term-size*:
assumes *match s t* = *Some* σ
and $x \in$ *vars-term t*
shows *size* (σ *x*) $\leq$ *size s*
using *assms vars-term-size* **by** (*metis match-matches*)

lemma [*termination-simp*]:
assumes *match* (*Fun f ss*) (*to-pterm l*) = *Some* σ
and $*$: (*s*, *t*) $\in$ *set* (*zip* (*map* σ (*vars-distinct l*)) *ts*)
shows *size s* $\leq$ *Suc* (*size-list size ss*)
proof −
from $*$ **have** *s* $\in$ *set* (*map* σ (*vars-distinct l*)) **by** (*blast elim*: *in-set-zipE*)

then obtain x **where** [*simp*]: $s = \sigma\ x$
and x: $x \in$ *vars-term* (*to-pterm* l) **by** (*induct* l) *auto*
from *match-vars-term-size* [*OF assms*(*1*) x]
show *?thesis* **by** *simp*
qed

Additional simp rule because we allow variable left-hand sides of rewrite rules at this point. Then *Var* x / α and α / *Var* x are also possible when evaluating residuals. This might become important when we want to introduce the error rule for residuals of composed proof terms.

lemma [*termination-simp*]:
assumes *match* (*Var* x) (*to-pterm* l) = *Some* σ
and $(a, b) \in$ *set* (*zip* (*map* σ (*vars-distinct* l)) *ts*)
shows *size* $a = 1$
proof−
from *assms*(*1*) **have** ∗:(*to-pterm* l) · σ = *Var* x **by** (*simp add*: *match-matches*)
then obtain y **where** y:l = *Var* y **by** (*metis subst-apply-eq-Var term.distinct*(*1*) *to-pterm.elims*)
with ∗ **have** ∗∗:$\sigma\ y$ = *Var* x **by** *simp*
from y **have** *vars-distinct* $l = [y]$ **by** (*simp add*: *vars-term-list.simps*(*1*))
with *assms*(*2*) y **have** a = *Var* x **by** (*simp add*: ∗∗ *in-set-zip*)
then show *?thesis* **by** *simp*
qed

```
fun residual :: ('f, 'v) pterm ⇒ ('f, 'v) pterm ⇒ ('f, 'v) pterm option (infixr re
70)
  where
  Var x re Var y =
    (if x = y then Some (Var x) else None)
| Pfun f As re Pfun g Bs =
    (if (f = g ∧ length As = length Bs) then
      (case those (map2 residual As Bs) of
        Some xs ⇒ Some (Pfun f xs)
      | None ⇒ None)
    else None)
| Prule α As re Prule β Bs =
    (if α = β then
      (case those (map2 residual As Bs) of
        Some xs ⇒ Some ((to-pterm (rhs α)) · ⟨xs⟩_α)
      | None ⇒ None)
    else None)
| Prule α As re B =
    (case match B (to-pterm (lhs α)) of
      None ⇒ None
    | Some σ ⇒
      (case those (map2 residual As (map σ (var-rule α))) of
        Some xs ⇒ Some (Prule α xs)
      | None ⇒ None))
| A re Prule α Bs =
```

```
  (case match A (to-pterm (lhs α)) of
    None ⇒ None
  | Some σ ⇒
    (case those (map2 residual (map σ (var-rule α)) Bs) of
      Some xs ⇒ Some ((to-pterm (rhs α)) · ⟨xs⟩α)
    | None ⇒ None))
| A re B = None
```

Since the interesting proofs about residuals always follow the same pattern of induction on the definition, we introduce the following 6 lemmas corresponding to the step cases.

```
lemma residual-fun-fun:
  assumes (Pfun f As) re (Pfun g Bs) = Some C
  shows f = g ∧ length As = length Bs ∧
      (∃ Cs. C = Pfun f Cs ∧
      length Cs = length As ∧
      (∀ i < length As. As!i re Bs!i = Some (Cs!i)))
proof−
  have ∗:f = g ∧ length As = length Bs
    using assms residual.simps(2) by (metis option.simps(3))
  then obtain Cs where Cs:those (map2 (re) As Bs) = Some Cs
    using assms residual.simps(2) option.simps(3) option.simps(4) by fastforce
  hence ∀ i < length As. As!i re Bs!i = Some (Cs!i)
    using ∗ those-some2 by fastforce
  with ∗ Cs assms(1) show ?thesis
    using length-those by fastforce
qed

lemma residual-rule-rule:
  assumes (Prule α As) re (Prule β Bs) = Some C
          (Prule α As) ∈ wf-pterm R
          (Prule β Bs) ∈ wf-pterm S
  shows α = β ∧ length As = length Bs ∧
      (∃ Cs. C = to-pterm (rhs α) · ⟨Cs⟩α ∧
      length Cs = length As ∧
      (∀ i < length As. As!i re Bs!i = Some (Cs!i)))
proof−
  have α = β
    using assms(1) residual.simps(3) by (metis option.simps(3))
  with assms(2,3) have l: length As = length Bs
    using length-args-well-Prule by blast
  from ‹α = β› obtain Cs where Cs:those (map2 (re) As Bs) = Some Cs
    using assms by fastforce
  hence ∀ i < length As. As!i re Bs!i = Some (Cs!i)
    using l those-some2 by fastforce
  with ‹α = β› l Cs assms(1) show ?thesis
    using length-those by fastforce
qed
```

lemma *residual-rule-var*:
assumes $(Prule\ \alpha\ As)\ re\ (Var\ x) = Some\ C$
$(Prule\ \alpha\ As) \in wf\text{-}pterm\ R$
shows $\exists\sigma.\ match\ (Var\ x)\ (to\text{-}pterm\ (lhs\ \alpha)) = Some\ \sigma\ \wedge$
$(\exists Cs.\ C = Prule\ \alpha\ Cs\ \wedge$
$length\ Cs = length\ As\ \wedge$
$(\forall i < length\ As.\ As!i\ re\ (\sigma\ (var\text{-}rule\ \alpha\ !\ i)) = Some\ (Cs!i)))$
proof−
from $assms(2)$ **have** l:$length\ As = length\ (var\text{-}rule\ \alpha)$
using *wf-pterm.simps* **by** *fastforce*
obtain σ **where** σ:$match\ (Var\ x)\ (to\text{-}pterm\ (lhs\ \alpha)) = Some\ \sigma$
using $assms(1)$ **by** *fastforce*
then obtain Cs **where** Cs:$those\ (map2\ residual\ As\ (map\ \sigma\ (var\text{-}rule\ \alpha))) = Some\ Cs$
using $assms(1)$ **by** *fastforce*
with l **have** $l2$:$length\ Cs = length\ As$
using *length-those* **by** *fastforce*
from Cs **have** $\forall i < length\ As.\ As!i\ re\ (\sigma\ (var\text{-}rule\ \alpha\ !\ i)) = Some\ (Cs!i)$
using *l those-some2* **by** *fastforce*
with $\sigma\ Cs\ assms(1)\ l2$ **show** *?thesis* **by** *simp*
qed

lemma *residual-rule-fun*:
assumes $(Prule\ \alpha\ As)\ re\ (Pfun\ f\ Bs) = Some\ C$
$(Prule\ \alpha\ As) \in wf\text{-}pterm\ R$
shows $\exists\sigma.\ match\ (Pfun\ f\ Bs)\ (to\text{-}pterm\ (lhs\ \alpha)) = Some\ \sigma\ \wedge$
$(\exists Cs.\ C = Prule\ \alpha\ Cs\ \wedge$
$length\ Cs = length\ As\ \wedge$
$(\forall i < length\ As.\ As!i\ re\ (\sigma\ (var\text{-}rule\ \alpha\ !\ i)) = Some\ (Cs!i)))$
proof−
from $assms(2)$ **have** l:$length\ As = length\ (var\text{-}rule\ \alpha)$
using *wf-pterm.simps* **by** *fastforce*
obtain σ **where** σ:$match\ (Pfun\ f\ Bs)\ (to\text{-}pterm\ (lhs\ \alpha)) = Some\ \sigma$
using $assms(1)$ **by** *fastforce*
then obtain Cs **where** Cs:$those\ (map2\ residual\ As\ (map\ \sigma\ (var\text{-}rule\ \alpha))) = Some\ Cs$
using $assms(1)$ **by** *fastforce*
with l **have** $l2$:$length\ Cs = length\ As$
using *length-those* **by** *fastforce*
from Cs **have** $\forall i < length\ As.\ As!i\ re\ (\sigma\ (var\text{-}rule\ \alpha\ !\ i)) = Some\ (Cs!i)$
using *l those-some2* **by** *fastforce*
with $\sigma\ Cs\ assms(1)\ l2$ **show** *?thesis* **by** *auto*
qed

lemma *residual-var-rule*:
assumes $(Var\ x)\ re\ (Prule\ \alpha\ As) = Some\ C$
$(Prule\ \alpha\ As) \in wf\text{-}pterm\ R$
shows $\exists\sigma.\ match\ (Var\ x)\ (to\text{-}pterm\ (lhs\ \alpha)) = Some\ \sigma\ \wedge$
$(\exists Cs.\ C = (to\text{-}pterm\ (rhs\ \alpha)) \cdot \langle Cs\rangle_\alpha\ \wedge$

length Cs = *length As* ∧
($\forall i < $ *length As*. (σ (*var-rule* α ! i) *re As*!i) = *Some* (*Cs*!i)))

proof−
from *assms*(2) **have** *l*:*length As* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
obtain σ **where** σ:*match* (*Var x*) (*to-pterm* (*lhs* α)) = *Some* σ
using *assms*(1) **by** *fastforce*
then obtain *Cs* **where** *Cs*:*those* (*map2 residual* (*map* σ (*var-rule* α)) *As*) = *Some Cs*
using *assms*(1) **by** *fastforce*
with *l* **have** *l2*:*length Cs* = *length As*
using *length-those* **by** *fastforce*
from *Cs* **have** $\forall i < $ *length As*. (σ (*var-rule* α ! i)) *re As*!i = *Some* (*Cs*!i)
using *l those-some2* **by** *fastforce*
with σ *Cs assms*(1) *l2* **show** *?thesis* **by** *auto*
qed

lemma *residual-fun-rule*:
assumes (*Pfun f Bs*) *re* (*Prule* α *As*) = *Some C*
(*Prule* α *As*) ∈ *wf-pterm R*
shows $\exists \sigma$. *match* (*Pfun f Bs*) (*to-pterm* (*lhs* α)) = *Some* σ ∧
($\exists$ *Cs*. *C* = (*to-pterm* (*rhs* α)) · $\langle Cs \rangle_\alpha$ ∧
length Cs = *length As* ∧
($\forall i < $ *length As*. (σ (*var-rule* α ! i)) *re As*!i = *Some* (*Cs*!i)))

proof−
from *assms*(2) **have** *l*:*length As* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
obtain σ **where** σ:*match* (*Pfun f Bs*) (*to-pterm* (*lhs* α)) = *Some* σ
using *assms*(1) **by** *fastforce*
then obtain *Cs* **where** *Cs*:*those* (*map2 residual* (*map* σ (*var-rule* α)) *As*) = *Some Cs*
using *assms*(1) **by** *fastforce*
with *l* **have** *l2*:*length Cs* = *length As*
using *length-those* **by** *fastforce*
with *Cs* **have** $\forall i < $ *length As*. (σ (*var-rule* α ! i)) *re As*!i = *Some* (*Cs*!i)
using *l those-some2* **by** *fastforce*
with σ *Cs assms*(1) *l2* **show** *?thesis* **by** *auto*
qed

t / A = *tgt*(A)

lemma *res-empty1*:
assumes *is-empty-step t co-initial A t A* ∈ *wf-pterm R*
shows *t re A* = *Some* (*to-pterm* (*target A*))
proof −
from *assms*(1,2) **have** *t* = *to-pterm* (*source A*)
by (*simp add*: *empty-coinitial*)
then show *?thesis* **using** *assms*(3) **proof** (*induction A arbitrary*: *t*)
case (*Var x*)
then show *?case* **by** *simp*

next
case (*Pfun f As*)
let *?ts* = (*map* (*to-pterm* ∘ *source*) *As*)
from *Pfun*(*3*) **have** ∀ *a* ∈ *set As*. *a* ∈ *wf-pterm R* **by** *blast*
with *Pfun*(*1*) **have** *those* (*map2 residual ?ts As*) = *Some* (*map* (*to-pterm* ∘ *target*) *As*) **by** (*simp add*:*those-some*)
then show *?case* **unfolding** *Pfun*(*2*) **by** *simp*
next
case (*Prule* α *As*)
let *?ts* = (*map* (*to-pterm* ∘ *source*) *As*)
from *Prule*(*3*) **have** *l*:*length ?ts* = *length* (*var-rule* α) **using** *wf-pterm.simps* **by** *fastforce*
moreover from *Prule*(*3*) **have** *well*:∀ *a* ∈ *set As*. *a* ∈ *wf-pterm R* **by** *blast*
from *Prule*(*1*) **have** *args*:*those* (*map2 residual ?ts As*) = *Some* (*map* (*to-pterm* ∘ *target*) *As*) **using** *well* **by** (*simp add*:*those-some*)
from *Prule*(*2*) **have** *t*:*t* = (*to-pterm* (*lhs* α)) · $\langle$*?ts*$\rangle_\alpha$ **by** (*simp add*: *to-pterm-subst*)

then obtain σ **where** σ:
match t (*to-pterm* (*lhs* α)) = *Some* σ
(∀ *x*∈ *set* (*var-rule* α). ($\langle$*?ts*$\rangle_\alpha$) *x* = σ *x*)
using *lhs-subst-trivial* **by** *blast*
from σ(*2*) *l* **have** *ts*:*map* σ (*var-rule* α) = *?ts* **by** (*smt apply-lhs-subst-var-rule map-eq-conv*)
from *Prule*(*1*) **have** *those* (*map2 residual ?ts As*) = *Some* (*map* (*to-pterm* ∘ *target*) *As*) **using** *well* **by** (*simp add*:*those-some*)
with *ts* **have** *args*:*those* (*map2 residual* (*map* σ (*var-rule* α)) *As*) = *Some* (*map* (*to-pterm* ∘ *target*) *As*) **by** *simp*
show *?case* **proof** (*cases t rule*:*source.cases*)
case (*1 x*)
then show *?thesis* **using** *args* σ(*1*) **by** (*simp add*: *to-pterm-subst*)
next
case (*2 f As*)
then show *?thesis* **using** *args* σ(*1*) **by** (*simp add*: *to-pterm-subst*)
next
case (*3* α *As*)
then show *?thesis* **using** *Prule*(*2*) **by** (*metis is-empty-step.simps*(*3*) *to-pterm-empty*)
qed
qed
qed

A / *t* = *A*

lemma *res-empty2*:
assumes *A* ∈ *wf-pterm R*
shows *A re* (*to-pterm* (*source A*)) = *Some A*
using *assms* **proof** (*induction A*)
case (*2 As f*)
then have *those* (*map2 residual As* (*map* (*to-pterm* ∘ *source*) *As*)) = *Some As* **by** (*simp add*:*those-some*)
then show *?case* **by** *simp*

next
  case (*3* α *As*)
  then have σ: *match* (*to-pterm* (*lhs* $\alpha \cdot \langle$*map source As*$\rangle_\alpha$)) (*to-pterm* (*lhs* α)) = *Some* ($\langle$*map* (*to-pterm* $\circ$ *source*) *As*$\rangle_\alpha$)
    by (*metis* (*no-types*, *lifting*) *fun-mk-subst lhs-subst-trivial map-map to-pterm.simps*(*1*) *to-pterm-subst*)
  from *3* **have** *those* (*map2 residual As* (*map* (*to-pterm* $\circ$ *source*) *As*)) = *Some As*
    by (*simp add*:*those-some*)
  then have *args*:*those* (*map2 residual As* (*map* ($\langle$*map* (*to-pterm* $\circ$ *source*) *As*$\rangle_\alpha$) (*var-rule* α))) = *Some As*
    by (*metis 3.hyps*(*2*) *apply-lhs-subst-var-rule length-map*)
  show *?case* **proof**(*cases source* (*Prule* α *As*))
    case (*Var x*)
    then show *?thesis*
      using σ *residual.simps*(*4*)[*of* α *As x*] *args* **by** *auto*
  next
    case (*Fun f ts*)
    then show *?thesis*
      using σ *residual.simps*(*5*)[*of* α *As f*] *args* **by** *auto*
  qed
qed *simp*

A / A = tgt(*A*)

lemma *res-same*: *A re A* = *Some* (*to-pterm* (*target A*))
proof(*induction A*)
case (*Var x*)
then show *?case* **by** *simp*
next
  case (*Pfun f As*)
  then have *list-all* ($\lambda x.\ x \neq$ *None*) (*map2 residual As As*) **by** (*simp add*: *list-all-length*)
  then obtain *xs* **where** *xs*:*those* (*map2 residual As As*) = *Some xs* **using** *those-not-none-xs* **by** *fastforce*
  then have *l*:*length As* = *length xs* **using** *length-those* **by** *fastforce*
  from *xs* **have** *IH*:*i* < *length As* $\Longrightarrow$ *xs*!*i* = *to-pterm* (*target* (*As*!*i*)) **for** *i* **using** *Pfun those-some2* **by** *force*
  from *IH l* **have** *map* (*to-pterm*$\circ$*target*) *As* = *xs* **by** (*simp add*: *map-nth-eq-conv*)
  then have *to-pterm* (*target* (*Pfun f As*)) = *Pfun f xs* **by** *simp*
  then show *?case* **using** *xs* **by** *simp*
next
  case (*Prule* α *As*)
  then have *list-all* ($\lambda x.\ x \neq$ *None*) (*map2 residual As As*) **by** (*simp add*: *list-all-length*)
  then obtain *xs* **where** *xs*:*those* (*map2 residual As As*) = *Some xs* **using** *those-not-none-xs* **by** *fastforce*
  then have *l*:*length As* = *length xs* **using** *length-those* **by** *fastforce*
  from *xs* **have** *IH*:*i* < *length As* $\Longrightarrow$ *xs*!*i* = *to-pterm* (*target* (*As*!*i*)) **for** *i* **using** *Prule those-some2* **by** *force*
  from *IH l* **have** $*$:*map* (*to-pterm*$\circ$*target*) *As* = *xs* **by** (*simp add*: *map-nth-eq-conv*)
  have *to-pterm* (*target* (*Prule* α *As*)) = *to-pterm* (*rhs* $\alpha \cdot \langle$*map target As*$\rangle_\alpha$) **by**

simp
  also have ... = (*to-pterm* (*rhs* α)) · (*to-pterm* ∘ ⟨*map target As*⟩$_\alpha$) **by** (*simp add*: *to-pterm-subst*)
  also have ... = (*to-pterm* (*rhs* α)) · ⟨*xs*⟩$_\alpha$ **using** ∗ **by** *simp*
  finally show *?case* **using** *xs* **by** *simp*
qed

lemma *residual-src-tgt*:
  assumes *A re B* = *Some C A* ∈ *wf-pterm R B* ∈ *wf-pterm S*
  shows *source C* = *target B*
  using *assms* **proof**(*induction A B arbitrary*: *C rule*: *residual.induct*)
  case (*1 x y*)
  then show *?case*
    by (*metis option.distinct*(*1*) *option.sel residual.simps*(*1*) *source.simps*(*1*) *target.simps*(*1*))
next
  case (*2 f As g Bs*)
  then obtain *Cs* **where** ∗:*f* = *g* ∧ *length As* = *length Bs* ∧
      C = *Pfun f Cs* ∧ *length Cs* = *length As* ∧
      (∀ *i*<*length As*. *As* ! *i re Bs* ! *i* = *Some* (*Cs* ! *i*))
    by (*meson residual-fun-fun*)
  then have *length* (*zip As Bs*) = *length As* **by** *simp*
  moreover from *2*(*3*) **have** ∀ *a* ∈ *set As*. *a* ∈ *wf-pterm R* **by** *blast*
  moreover from *2*(*4*) **have** ∀ *b* ∈ *set Bs*. *b* ∈ *wf-pterm S* **by** *blast*
  ultimately have ∀ *i* < *length As*. *source* (*Cs*!*i*) = *target* (*Bs*!*i*)
    using ∗ *2* **by** (*metis nth-mem nth-zip*)
  with ∗ **show** *?case* **by** (*simp add*: *nth-map-conv*)
next
  case (*3* α *As* β *Bs*)
  then obtain *Cs* **where** ∗:α = β ∧ *length As* = *length Bs* ∧
      C = *to-pterm* (*rhs* α) · ⟨*Cs*⟩$_\alpha$ ∧ *length Cs* = *length As* ∧
      (∀ *i*<*length As*. *As* ! *i re Bs* ! *i* = *Some* (*Cs* ! *i*))
    by (*meson residual-rule-rule*)
  then have *length* (*zip As Bs*) = *length As* **by** *simp*
  moreover from *3*(*3*) **have** ∀ *a* ∈ *set As*. *a* ∈ *wf-pterm R* **by** *blast*
  moreover from *3*(*4*) **have** ∀ *b* ∈ *set Bs*. *b* ∈ *wf-pterm S* **by** *blast*
  ultimately have *IH*:∀ *i* < *length As*. *source* (*Cs*!*i*) = *target* (*Bs*!*i*)
    using ∗ *3* **by** (*metis nth-mem nth-zip*)
  from ∗ **have** *source C* = (*rhs* β) · ⟨*map source Cs*⟩$_\beta$
     by (*simp add*: *source-apply-subst*)
  also have ... = (*rhs* β) · ⟨*map target Bs*⟩$_\beta$
    using ∗ *IH* **by** (*metis nth-map-conv*)
  finally show *?case* **by** *simp*
next
  case (*4-1* α *As v*)
  then obtain *Cs* σ **where** ∗:*match* (*Var v*) (*to-pterm* (*lhs* α)) = *Some* σ ∧
        C = *Prule* α *Cs* ∧ *length Cs* = *length As* ∧
        (∀ *i*<*length As*. *As* ! *i re* σ (*var-rule* α ! *i*) = *Some* (*Cs* ! *i*))
    by (*meson residual-rule-var*)

then obtain *Bs* **where** *Bs*:*length Bs* = *length* (*var-rule* α) $\wedge$
Var v = (*to-pterm* (*lhs* α)) $\cdot$ $\langle Bs \rangle_\alpha$ $\wedge$
($\forall$ *x* $\in$ *set* (*var-rule* α). σ *x* = ($\langle Bs \rangle_\alpha$) *x*)
using *match-lhs-subst* **by** *blast*
from *4-1*(*3*) **have** *as*:$\forall$ *a* $\in$ *set As*. *a* $\in$ *wf-pterm R* **by** *blast*
from *4-1*(*3*) **have** *l*:*length As* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
with $*$ **have** *well*:$\forall$ *i* < *length As*. σ (*var-rule* α ! *i*) $\in$ *wf-pterm S*
using *4-1*(*4*) **by** (*metis match-well-def vars-to-pterm*)
from *l* **have** *length* (*zip As* (*map* σ (*var-rule* α))) = *length As* **by** *simp*
with *4-1*(*1*,*3*) *well* $*$ *l as* **have** *IH*:$\forall$ *i* < *length As*. *source* (*Cs*!*i*) = *target* (*map* ($\langle Bs \rangle_\alpha$) (*var-rule* α) !*i*)
using *Bs* **by** (*smt length-map nth-map nth-mem nth-zip*)
from $*$ **have** *source C* = (*lhs* α) $\cdot$ $\langle$*map source Cs*$\rangle_\alpha$
by (*simp add*: *source-apply-subst*)
also have ... = (*lhs* α) $\cdot$ $\langle$*map* (*target* $\circ$ ($\langle Bs \rangle_\alpha$)) (*var-rule* α)$\rangle_\alpha$
using $*$ *l IH* **by** (*smt map-eq-conv′ map-map*)
also have ... = (*lhs* α) $\cdot$ (*target* $\circ$ ($\langle Bs \rangle_\alpha$))
using *Bs* **by** (*metis* (*no-types*, *lifting*) *apply-lhs-subst-var-rule fun-mk-subst map-map target.simps*(*1*))
also have ... = *target* (*to-pterm* (*lhs* α) $\cdot$ $\langle Bs \rangle_\alpha$)
by (*metis target-empty-apply-subst target-to-pterm to-pterm-empty*)
finally show *?case*
using *Bs* **by** *fastforce*
next
case (*4-2* α *As f Bs*)
then obtain *Cs* σ **where** $*$:*match* (*Pfun f Bs*) (*to-pterm* (*lhs* α)) = *Some* σ $\wedge$
C = *Prule* α *Cs* $\wedge$ *length Cs* = *length As* $\wedge$
($\forall$ *i*<*length As*. *As* ! *i re* σ (*var-rule* α ! *i*) = *Some* (*Cs* ! *i*))
by (*meson residual-rule-fun*)
then obtain *Bs′* **where** *Bs′*:*length Bs′* = *length* (*var-rule* α) $\wedge$
Pfun f Bs = (*to-pterm* (*lhs* α)) $\cdot$ $\langle Bs' \rangle_\alpha$ $\wedge$
($\forall$ *x* $\in$ *set* (*var-rule* α). σ *x* = ($\langle Bs' \rangle_\alpha$) *x*)
using *match-lhs-subst* **by** *blast*
from *4-2*(*3*) **have** *as*:$\forall$ *a* $\in$ *set As*. *a* $\in$ *wf-pterm R* **by** *blast*
from *4-2*(*3*) **have** *l*:*length As* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
with $*$ **have** *well*:$\forall$ *i* < *length As*. σ (*var-rule* α ! *i*) $\in$ *wf-pterm S*
using *4-2*(*4*) **by** (*metis match-well-def vars-to-pterm*)
from *l* **have** *length* (*zip As* (*map* σ (*var-rule* α))) = *length As* **by** *simp*
with *4-2*(*1*,*3*) *well* $*$ *l as* **have** *IH*:$\forall$ *i* < *length As*. *source* (*Cs*!*i*) = *target* (*map* ($\langle Bs' \rangle_\alpha$) (*var-rule* α) !*i*)
using *Bs′* **by** (*smt length-map nth-map nth-mem nth-zip*)
from $*$ **have** *source C* = (*lhs* α) $\cdot$ $\langle$*map source Cs*$\rangle_\alpha$
by (*simp add*: *source-apply-subst*)
also have ... = (*lhs* α) $\cdot$ $\langle$*map* (*target* $\circ$ ($\langle Bs' \rangle_\alpha$)) (*var-rule* α)$\rangle_\alpha$
using $*$ *l IH* **by** (*smt map-eq-conv′ map-map*)
also have ... = (*lhs* α) $\cdot$ (*target* $\circ$ ($\langle Bs' \rangle_\alpha$))
using *Bs′* **by** (*metis* (*no-types*, *lifting*) *apply-lhs-subst-var-rule fun-mk-subst*

map-map target.simps(*1*))
also have ... = *target* (*to-pterm* (*lhs* α) $\cdot$ $\langle Bs' \rangle_\alpha$)
by (*metis target-empty-apply-subst target-to-pterm to-pterm-empty*)
finally show *?case*
using *Bs'* **by** *fastforce*
next
case (*5-1 v* α *As*)
then obtain *Cs* σ **where** $*$:*match* (*Var v*) (*to-pterm* (*lhs* α)) = *Some* σ $\wedge$
C = *to-pterm* (*rhs* α) $\cdot$ $\langle Cs \rangle_\alpha$ $\wedge$ *length Cs* = *length As* $\wedge$
($\forall$ *i*<*length As*. σ (*var-rule* α ! *i*) *re As* ! *i* = *Some* (*Cs* ! *i*))
by (*meson residual-var-rule*)
from *5-1*(*4*) **have** *as*:$\forall$ *a* $\in$ *set As*. *a* $\in$ *wf-pterm S* **by** *blast*
from *5-1*(*4*) **have** *l*:*length As* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
with $*$ **have** *well*:$\forall$ *i* < *length As*. σ (*var-rule* α ! *i*) $\in$ *wf-pterm R*
using *5-1*(*3*) **by** (*metis match-well-def vars-to-pterm*)
from *l* **have** *length* (*zip* (*map* σ (*var-rule* α)) *As*) = *length As* **by** *simp*
with *5-1*(*1*,*4*) *well* $*$ *l as* **have** *IH*:$\forall$ *i* < *length As*. *source* (*Cs*!*i*) = *target* (*As*!*i*)

by (*smt length-map nth-map nth-mem nth-zip*)
from $*$ **have** *source C* = (*rhs* α) $\cdot$ $\langle$*map source Cs*$\rangle_\alpha$
by (*simp add*: *source-apply-subst*)
also have ... = (*rhs* α) $\cdot$ $\langle$*map target As*$\rangle_\alpha$
using $*$ *IH* **by** (*metis* (*no-types*, *lifting*) *map-eq-conv'*)
finally show *?case* **by** *simp*
next
case (*5-2 f Bs* α *As*)
then obtain *Cs* σ **where** $*$:*match* (*Pfun f Bs*) (*to-pterm* (*lhs* α)) = *Some* σ $\wedge$
C = *to-pterm* (*rhs* α) $\cdot$ $\langle Cs \rangle_\alpha$ $\wedge$ *length Cs* = *length As* $\wedge$
($\forall$ *i*<*length As*. σ (*var-rule* α ! *i*) *re As* ! *i* = *Some* (*Cs* ! *i*))
by (*meson residual-fun-rule*)
from *5-2*(*4*) **have** *as*:$\forall$ *a* $\in$ *set As*. *a* $\in$ *wf-pterm S* **by** *blast*
from *5-2*(*4*) **have** *l*:*length As* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
with $*$ **have** *well*:$\forall$ *i* < *length As*. σ (*var-rule* α ! *i*) $\in$ *wf-pterm R*
using *5-2*(*3*) **by** (*metis match-well-def vars-to-pterm*)
from *l* **have** *length* (*zip* (*map* σ (*var-rule* α)) *As*) = *length As* **by** *simp*
with *5-2*(*1*,*4*) *well* $*$ *l as* **have** *IH*:$\forall$ *i* < *length As*. *source* (*Cs*!*i*) = *target* (*As*!*i*)

by (*smt length-map nth-map nth-mem nth-zip*)
from $*$ **have** *source C* = (*rhs* α) $\cdot$ $\langle$*map source Cs*$\rangle_\alpha$
by (*simp add*: *source-apply-subst*)
also have ... = (*rhs* α) $\cdot$ $\langle$*map target As*$\rangle_\alpha$
using $*$ *IH* **by** (*metis* (*no-types*, *lifting*) *map-eq-conv'*)
finally show *?case* **by** *simp*
qed *simp-all*

The following two lemmas are used inside the induction proof for the result $tgt(A\ /\ B) = tgt(B\ /\ A)$. Defining them here, outside the main proof makes them reusable for the symmetric cases of the proof.

lemma *tgt-tgt-rule-var*:
  assumes $\bigwedge\sigma$ *a b c d. match* (*Var v*) (*to-pterm* (*lhs* α)) = *Some* $\sigma \Longrightarrow$
        (*a*,*b*) $\in$ *set* (*zip As* (*map* σ (*var-rule* α))) $\Longrightarrow$
          a re b = *Some c* $\Longrightarrow$ *b re a* = *Some d* $\Longrightarrow$ *a* $\in$ *wf-pterm R* $\Longrightarrow$ *b* $\in$ *wf-pterm S* $\Longrightarrow$
          target c = *target d*
        Prule α *As re Var v* = *Some C*
        Var v re Prule α *As* = *Some D*
        Prule α *As* $\in$ *wf-pterm R*
  shows *target C* = *target D*
proof−
  from *assms*(*4*) **have** *l*:*length As* = *length* (*var-rule* α)
    using *wf-pterm.simps* **by** *fastforce*
  from *assms*(*4*) **have** *as*:$\forall$ *a* $\in$ *set As. a* $\in$ *wf-pterm R* **by** *blast*
  from *assms*(*2*,*4*) **obtain** σ **where** σ:*match* (*Var v*) (*to-pterm* (*lhs* α)) = *Some* σ
    by (*meson residual-rule-var*)
  with *l* **have** *well-def*:$\forall$ *i* < *length As.* σ (*var-rule* α ! *i*) $\in$ *wf-pterm S*
    using *match-well-def* **by** (*metis vars-to-pterm wf-pterm.intros*(*1*))
  from *assms*(*2*,*4*) σ **obtain** *Cs* **where** *Cs*:
      C = *Prule* α *Cs* $\wedge$ *length Cs* = *length As*
      ($\forall$ *i*<*length As. As* ! *i re* σ (*var-rule* α ! *i*) = *Some* (*Cs* ! *i*))
    by (*metis option.inject residual-rule-var*)
  from *assms*(*3*,*4*) σ **obtain** *Ds* **where** *Ds*:
      D = *to-pterm* (*rhs* α) $\cdot$ $\langle Ds\rangle_\alpha$ $\wedge$ *length Ds* = *length As*
      ($\forall$ *i*<*length As.* σ (*var-rule* α ! *i*) *re As* ! *i* = *Some* (*Ds* ! *i*))
    by (*metis option.inject residual-var-rule*)
  from *l* **have** *length As* = *length* (*zip As* (*map* σ (*var-rule* α)))
    by *simp*
  with *assms*(*1*,*4*) σ *l Cs*(*2*) *Ds*(*2*) *well-def* **have** *IH*:$\forall$ *i* < *length As. target* (*Cs*!*i*) = *target* (*Ds*!*i*)
    using *as* **by** (*smt length-map nth-map nth-mem nth-zip*)
  from *Cs* **have** *target C* = (*rhs* α) $\cdot$ $\langle$*map target Cs*$\rangle_\alpha$ **by** *simp*
  moreover from *Ds*(*1*) **have** *target D* = (*rhs* α) $\cdot$ $\langle$*map target Ds*$\rangle_\alpha$
    using *target-empty-apply-subst to-pterm-empty* **by** (*metis fun-mk-subst target.simps*(*1*) *target-to-pterm*)
  ultimately show *?thesis*
    using *IH Cs*(*1*) *Ds*(*1*) **by** (*metis nth-map-conv*)
qed

lemma *tgt-tgt-rule-fun*:
  assumes $\bigwedge\sigma$ *a b c d. match* (*Pfun f Bs*) (*to-pterm* (*lhs* α)) = *Some* $\sigma \Longrightarrow$
        (*a*,*b*) $\in$ *set* (*zip As* (*map* σ (*var-rule* α))) $\Longrightarrow$
          a re b = *Some c* $\Longrightarrow$ *b re a* = *Some d* $\Longrightarrow$ *a* $\in$ *wf-pterm R* $\Longrightarrow$ *b* $\in$ *wf-pterm S* $\Longrightarrow$
          target c = *target d*
        Prule α *As re Pfun f Bs* = *Some C*
        Pfun f Bs re Prule α *As* = *Some D*

Prule α As ∈ wf-pterm R
Pfun f Bs ∈ wf-pterm S
shows *target C = target D*
proof −
from *assms(4)* **have** *l:length As = length (var-rule α)*
using *wf-pterm.simps* **by** *fastforce*
from *assms(4)* **have** *as:∀ a ∈ set As. a ∈ wf-pterm R* **by** *blast*
from *assms(2,4)* **obtain** *σ* **where** *σ:match (Pfun f Bs) (to-pterm (lhs α)) = Some σ*
by *(meson residual-rule-fun)*
with *assms(2,5) l* **have** *well-def:∀ i < length As. σ (var-rule α ! i) ∈ wf-pterm S*
using *match-well-def* **by** *(metis vars-to-pterm)*
from *assms(2,4) σ* **obtain** *Cs* **where** *Cs*:
C = Prule α Cs ∧ length Cs = length As
(∀ i<length As. As ! i re σ (var-rule α ! i) = Some (Cs ! i))
by *(metis option.inject residual-rule-fun)*
from *assms(3,4) σ* **obtain** *Ds* **where** *Ds*:
D = to-pterm (rhs α) · ⟨Ds⟩$_\alpha$ *∧ length Ds = length As*
(∀ i<length As. σ (var-rule α ! i) re As ! i = Some (Ds ! i))
by *(metis option.inject residual-fun-rule)*
from *l* **have** *length As = length (zip As (map σ (var-rule α)))*
by *simp*
with *assms(1,4,5) σ l Cs(2) Ds(2) well-def* **have** *IH:∀ i < length As. target (Cs!i) = target (Ds!i)*
using *as* **by** *(smt length-map nth-map nth-mem nth-zip)*
from *Cs* **have** *target C = (rhs α) · ⟨map target Cs⟩*$_\alpha$ **by** *simp*
moreover from *Ds(1)* **have** *target D = (rhs α) · ⟨map target Ds⟩*$_\alpha$
using *target-empty-apply-subst to-pterm-empty* **by** *(metis fun-mk-subst target.simps(1) target-to-pterm)*
ultimately show *?thesis*
using *IH Cs(1) Ds(1)* **by** *(metis nth-map-conv)*
qed

lemma *residual-tgt-tgt*:
assumes *A re B = Some C B re A = Some D A ∈ wf-pterm R B ∈ wf-pterm S*
shows *target C = target D*
using *assms* **proof**(*induction A B arbitrary: C D rule:residual.induct*)
case *(1 x y)*
then show *?case* **by** *(metis option.sel residual.simps(1))*
next
case *(2 f As g Bs)*
from *2(4)* **have** *as:∀ a ∈ set As. a ∈ wf-pterm R* **by** *blast*
from *2(5)* **have** *bs:∀ b ∈ set Bs. b ∈ wf-pterm S* **by** *blast*
let *?l = length As*
from *2(2)* **have** *∗: f = g ∧ ?l = length Bs*
by *(meson residual-fun-fun)*
from *2(2)* **obtain** *Cs* **where** *Cs*:
C = Pfun f Cs ∧ length Cs = ?l ∧ (∀ i < ?l. As ! i re Bs ! i = Some (Cs ! i))

by (*meson residual-fun-fun*)
from *2*(*3*) **obtain** *Ds* **where** *Ds*:
$D = Pfun\ g\ Ds \wedge length\ Ds = ?l \wedge (\forall i < ?l.\ Bs\ !\ i\ re\ As\ !\ i = Some\ (Ds\ !\ i))$
using ∗ **by** (*metis residual-fun-fun*)
from ∗ **have** $length\ (zip\ As\ Bs) = ?l$ **by** *simp*
with *2*(*1*,*4*,*5*) ∗ *Cs Ds* **have** $\forall i < ?l.\ target\ (Cs!i) = target\ (Ds!i)$
using *as bs* **by** (*metis nth-mem nth-zip*)
with ∗ *Cs Ds* **show** *?case*
by (*simp add*: *map-nth-eq-conv*)
next
case (*3* α *As* β *Bs*)
from *3*(*4*) **have** *as*:$\forall a \in set\ As.\ a \in wf\text{-}pterm\ R$ **by** *blast*
from *3*(*5*) **have** *bs*:$\forall b \in set\ Bs.\ b \in wf\text{-}pterm\ S$ **by** *blast*
let $?l = length\ As$
from *3*(*2*,*4*,*5*) **have** ∗: $\alpha = \beta \wedge ?l = length\ Bs$
by (*meson residual-rule-rule*)
from *3*(*2*,*4*,*5*) **obtain** *Cs* **where** *Cs*:
$C = to\text{-}pterm\ (rhs\ \alpha) \cdot \langle Cs\rangle_\alpha \wedge length\ Cs = ?l \wedge (\forall i < ?l.\ As\ !\ i\ re\ Bs\ !\ i = Some\ (Cs\ !\ i))$
by (*meson residual-rule-rule*)
from *3*(*3*,*4*,*5*) **obtain** *Ds* **where** *Ds*:
$D = to\text{-}pterm\ (rhs\ \alpha) \cdot \langle Ds\rangle_\alpha \wedge length\ Ds = ?l \wedge (\forall i < ?l.\ Bs\ !\ i\ re\ As\ !\ i = Some\ (Ds\ !\ i))$
using ∗ **by** (*metis residual-rule-rule*)
from ∗ **have** $length\ (zip\ As\ Bs) = ?l$ **by** *simp*
with *3*(*1*,*4*,*5*) ∗ *Cs Ds* **have** *IH*:$\forall i < ?l.\ target\ (Cs!i) = target\ (Ds!i)$
using *as bs* **by** (*metis nth-mem nth-zip*)
from *Cs* **have** $target\ C = (rhs\ \alpha) \cdot \langle map\ target\ Cs\rangle_\alpha$
using *target-empty-apply-subst to-pterm-empty* **by** (*metis fun-mk-subst target.simps*(*1*) *target-to-pterm*)
moreover from *Ds* **have** $target\ D = (rhs\ \alpha) \cdot \langle map\ target\ Ds\rangle_\alpha$
using *target-empty-apply-subst to-pterm-empty* **by** (*metis fun-mk-subst target.simps*(*1*) *target-to-pterm*)
ultimately show *?case*
using *IH Cs Ds* **by** (*metis nth-map-conv*)
next
case (*4-1* α *As v*)
then show *?case* **using** *tgt-tgt-rule-var* **by** *fastforce*
next
case (*4-2* α *As f Bs*)
then show *?case* **using** *tgt-tgt-rule-fun* **by** *fastforce*
next
case (*5-1 v* α *As*)
from *5-1*(*1*) **have** $match\ (Var\ v)\ (to\text{-}pterm\ (lhs\ \alpha)) = Some\ \sigma \Longrightarrow$
$(a,b) \in set\ (zip\ As\ (map\ \sigma\ (var\text{-}rule\ \alpha))) \Longrightarrow$
$a\ re\ b = Some\ c \Longrightarrow b\ re\ a = Some\ d \Longrightarrow a \in wf\text{-}pterm\ S \Longrightarrow b \in wf\text{-}pterm\ R \Longrightarrow$
$target\ c = target\ d$ **for** $\sigma\ a\ b\ c\ d$
using *zip-symm* **by** *fastforce*

with *5-1*(*2*,*3*,*5*) **have** *target D* = *target C* **using** *tgt-tgt-rule-var* **by** *fastforce*
then show *?case* **by** *simp*
next
case (*5-2 f Bs* α *As*)
from *5-2*(*1*) **have** *match* (*Pfun f Bs*) (*to-pterm* (*lhs* α)) = *Some* $\sigma \Longrightarrow$
(*a*,*b*) $\in$ *set* (*zip As* (*map* σ (*var-rule* α))) $\Longrightarrow$
a re b = *Some c* $\Longrightarrow$ *b re a* = *Some d* $\Longrightarrow$ *a* $\in$ *wf-pterm S* $\Longrightarrow$ *b* $\in$ *wf-pterm R* $\Longrightarrow$
target c = *target d* **for** σ *a b c d*
using *zip-symm* **by** *fastforce*
with *5-2*(*2*,*3*,*4*,*5*) **have** *target D* = *target C* **using** *tgt-tgt-rule-fun* **by** *fastforce*
then show *?case* **by** *simp*
qed *simp-all*

lemma *rule-residual-lhs*:
assumes *args*:*those* (*map2* (*re*) *As Bs*) = *Some Cs*
and *is-Fun*:*is-Fun* (*lhs* α) **and** *l*:*length Bs* = *length* (*var-rule* α)
shows *Prule* α *As re* (*to-pterm* (*lhs* α) $\cdot$ $\langle Bs \rangle_\alpha$) = *Some* (*Prule* α *Cs*)
proof−
from *is-Fun* **obtain** *f ts* **where** *lhs* α = *Fun f ts*
by *auto*
then have *f*:*to-pterm* (*lhs* α) $\cdot$ $\langle Bs \rangle_\alpha$ = *Pfun f* (*map* ($\lambda t.\ t \cdot \langle Bs \rangle_\alpha$) (*map to-pterm ts*))
by *simp*
then have *match*:*match* ((*to-pterm* (*lhs* α)) $\cdot$ $\langle Bs \rangle_\alpha$) (*to-pterm* (*lhs* α)) = *Some* ($\langle Bs \rangle_\alpha$)
using *lhs-subst-trivial* **by** *blast*
from *l* **have** *map* ($\langle Bs \rangle_\alpha$) (*var-rule* α) = *Bs*
using *apply-lhs-subst-var-rule* **by** *blast*
with *args* **have** *those* (*map2* (*re*) *As* (*map* ($\langle Bs \rangle_\alpha$) (*var-rule* α))) = *Some Cs*
by *presburger*
then show *?thesis*
using *residual.simps*(*5*) *match* **unfolding** *f* **by** *auto*
qed

lemma *residual-well-defined*:
assumes *A* $\in$ *wf-pterm R B* $\in$ *wf-pterm S A re B* = *Some C*
shows *C* $\in$ *wf-pterm R*
using *assms* **proof**(*induction A B arbitrary*:*C rule*:*residual.induct*)
case (*1 x y*)
then show *?case*
by (*metis option.distinct*(*1*) *option.sel residual.simps*(*1*))
next
case (*2 f As g Bs*)
then obtain *Cs* **where** *f* = *g* $\wedge$ *length As* = *length Bs* $\wedge$
C = *Pfun f Cs* $\wedge$
length Cs = *length As* $\wedge$
($\forall$ *i* < *length As. As*!*i re Bs*!*i* = *Some* (*Cs*!*i*))
by (*meson residual-fun-fun*)

moreover with *2* **have** $i < length\ As \Longrightarrow Cs!i \in wf\text{-}pterm\ R$ **for** *i*
using *fun-well-arg* **by** (*metis* (*no-types*, *opaque-lifting*) *fst-conv in-set-zip nth-mem snd-conv*)
ultimately show *?case*
by (*metis in-set-conv-nth wf-pterm.intros*(*2*))
next
case (*3* α *As* β *Bs*)
then obtain *Cs* **where** $\alpha = \beta \wedge length\ As = length\ Bs \wedge$
$(C = to\text{-}pterm\ (rhs\ \alpha) \cdot \langle Cs\rangle_\alpha \wedge$
$length\ Cs = length\ As \wedge$
$(\forall\, i < length\ As.\ As!i\ re\ Bs!i = Some\ (Cs!i)))$
by (*meson residual-rule-rule*)
moreover with *3* **have** $i < length\ As \Longrightarrow Cs!i \in wf\text{-}pterm\ R$ **for** *i*
using *fun-well-arg* **by** (*metis* (*no-types*, *opaque-lifting*) *fst-conv in-set-zip nth-mem snd-conv*)
ultimately show *?case*
by (*metis to-pterm-wf-pterm lhs-subst-well-def*)
next
case (*4-1* α *As v*)
then obtain *Cs* σ **where** $*{:}match\ (Var\ v)\ (to\text{-}pterm\ (lhs\ \alpha)) = Some\ \sigma \wedge$
$C = Prule\ \alpha\ Cs \wedge$
$length\ Cs = length\ As \wedge$
$(\forall\, i < length\ As.\ As!i\ re\ (\sigma\ (var\text{-}rule\ \alpha\ !\ i)) = Some\ (Cs!i))$
by (*meson residual-rule-var*)
from *4-1*(*2*) **have** $wellA{:}\forall\, i < length\ As.\ As!i \in wf\text{-}pterm\ R$
by *auto*
from *4-1*(*2*) **have** $l{:}\ length\ As = length\ (var\text{-}rule\ \alpha)$
using *wf-pterm.simps* **by** *fastforce*
then have $l2{:}length\ As = length\ (zip\ As\ (map\ \sigma\ (var\text{-}rule\ \alpha)))$
by *simp*
from *l* * **have** $\forall\, i < length\ As.\ \sigma\ (var\text{-}rule\ \alpha\ !\ i) \in wf\text{-}pterm\ S$
using *4-1*(*3*) **by** (*metis match-well-def vars-to-pterm*)
with *4-1*(*1*) * *wellA l2* **have** $\forall\, i < length\ As.\ Cs!i \in wf\text{-}pterm\ R$
by (*smt* (*z3*) *l length-map nth-map nth-mem nth-zip*)
with *4-1*(*2*) * **show** *?case*
by (*smt* (*verit*) *Inr-not-Inl in-set-conv-nth term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.cases wf-pterm.intros*(*3*))
next
case (*4-2* α *As f Bs*)
then obtain σ *Cs* **where** $*{:}match\ (Pfun\ f\ Bs)\ (to\text{-}pterm\ (lhs\ \alpha)) = Some\ \sigma \wedge$
$(C = Prule\ \alpha\ Cs \wedge$
$length\ Cs = length\ As \wedge$
$(\forall\, i < length\ As.\ As!i\ re\ (\sigma\ (var\text{-}rule\ \alpha\ !\ i)) = Some\ (Cs!i)))$
by (*meson residual-rule-fun*)
from *4-2*(*2*) **have** $wellA{:}\forall\, i < length\ As.\ As!i \in wf\text{-}pterm\ R$
by *auto*
from *4-2*(*2*) **have** $l{:}\ length\ As = length\ (var\text{-}rule\ \alpha)$
using *wf-pterm.simps* **by** *fastforce*
then have $l2{:}length\ As = length\ (zip\ As\ (map\ \sigma\ (var\text{-}rule\ \alpha)))$

    by *simp*
  from *l* $*$ **have** $\forall\, i < \mathit{length}\ \mathit{As}.\ \sigma\ (\mathit{var\text{-}rule}\ \alpha\ !\ i) \in \mathit{wf\text{-}pterm}\ S$
    using *4-2(3)* **by** (*metis match-well-def vars-to-pterm*)
  with *4-2(1)* $*$ *wellA l2* **have** $\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{Cs}!i \in \mathit{wf\text{-}pterm}\ R$
    by (*smt l length-map nth-map nth-mem nth-zip*)
  with *4-2(2)* $*$ **show** *?case*
      by (*smt* (*verit*, *ccfv-threshold*) *Inr-not-Inl in-set-conv-nth term.distinct(1)* *term.inject(2) wf-pterm.cases wf-pterm.intros(3)*)
next
  case (*5-1 v* α *As*)
  then obtain *Cs* σ **where** $*$:$\mathit{match}\ (\mathit{Var}\ v)\ (\mathit{to\text{-}pterm}\ (\mathit{lhs}\ \alpha)) = \mathit{Some}\ \sigma\ \wedge$
      $C = \mathit{to\text{-}pterm}\ (\mathit{rhs}\ \alpha) \cdot \langle \mathit{Cs}\rangle_\alpha\ \wedge$
      $\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As}\ \wedge$
      $(\forall\, i < \mathit{length}\ \mathit{As}.\ (\sigma\ (\mathit{var\text{-}rule}\ \alpha\ !\ i))\ \mathit{re}\ \mathit{As}!i = \mathit{Some}\ (\mathit{Cs}!i))$
    by (*meson residual-var-rule*)
  from *5-1(3)* **have** *wellA*:$\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \in \mathit{wf\text{-}pterm}\ S$
    by *auto*
  from *5-1(3)* **have** *l*: $\mathit{length}\ \mathit{As} = \mathit{length}\ (\mathit{var\text{-}rule}\ \alpha)$
    using *wf-pterm.simps* **by** *fastforce*
  then have *l2*:$\mathit{length}\ \mathit{As} = \mathit{length}\ (\mathit{zip}\ (\mathit{map}\ \sigma\ (\mathit{var\text{-}rule}\ \alpha))\ \mathit{As})$
    by *simp*
  from *l* $*$ **have** $\forall\, i < \mathit{length}\ \mathit{As}.\ \sigma\ (\mathit{var\text{-}rule}\ \alpha\ !\ i) \in \mathit{wf\text{-}pterm}\ R$
    using *5-1(2)* **by** (*metis match-well-def vars-to-pterm*)
  with *5-1(1)* $*$ *wellA l2* **have** $\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{Cs}!i \in \mathit{wf\text{-}pterm}\ R$
    by (*smt l length-map nth-map nth-mem nth-zip*)
  with $*$ **show** *?case*
    by (*metis lhs-subst-well-def to-pterm-wf-pterm*)
next
  case (*5-2 f Bs* α *As*)
  then obtain *Cs* σ **where** $*$:$\mathit{match}\ (\mathit{Pfun}\ f\ \mathit{Bs})\ (\mathit{to\text{-}pterm}\ (\mathit{lhs}\ \alpha)) = \mathit{Some}\ \sigma\ \wedge$
      $C = \mathit{to\text{-}pterm}\ (\mathit{rhs}\ \alpha) \cdot \langle \mathit{Cs}\rangle_\alpha\ \wedge$
      $\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As}\ \wedge$
      $(\forall\, i < \mathit{length}\ \mathit{As}.\ (\sigma\ (\mathit{var\text{-}rule}\ \alpha\ !\ i))\ \mathit{re}\ \mathit{As}!i = \mathit{Some}\ (\mathit{Cs}!i))$
    by (*meson residual-fun-rule*)
  from *5-2(3)* **have** *wellA*:$\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \in \mathit{wf\text{-}pterm}\ S$
    by *auto*
  from *5-2(3)* **have** *l*: $\mathit{length}\ \mathit{As} = \mathit{length}\ (\mathit{var\text{-}rule}\ \alpha)$
    using *wf-pterm.simps* **by** *fastforce*
  then have *l2*:$\mathit{length}\ \mathit{As} = \mathit{length}\ (\mathit{zip}\ (\mathit{map}\ \sigma\ (\mathit{var\text{-}rule}\ \alpha))\ \mathit{As})$
    by *simp*
  from *l* $*$ **have** $\forall\, i < \mathit{length}\ \mathit{As}.\ \sigma\ (\mathit{var\text{-}rule}\ \alpha\ !\ i) \in \mathit{wf\text{-}pterm}\ R$
    using *5-2(2)* **by** (*metis match-well-def vars-to-pterm*)
  with *5-2(1)* $*$ *wellA l2* **have** $\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{Cs}!i \in \mathit{wf\text{-}pterm}\ R$
    by (*smt l length-map nth-map nth-mem nth-zip*)
  with $*$ **show** *?case*
    by (*metis lhs-subst-well-def to-pterm-wf-pterm*)
qed *simp-all*

no-notation *sup* (**infixl** $\sqcup$ *65*)

3.2 Join

```
fun join :: ('f, 'v) pterm ⇒ ('f,'v) pterm ⇒ ('f,'v) pterm option (infixr ⊔ 70)
  where
  Var x ⊔ Var y =
    (if x = y then Some (Var x) else None)
| Pfun f As ⊔ Pfun g Bs =
    (if (f = g ∧ length As = length Bs) then
      (case those (map2 (⊔) As Bs) of
        Some xs ⇒ Some (Pfun f xs)
      | None ⇒ None)
    else None)
| Prule α As ⊔ Prule β Bs =
    (if α = β then
      (case those (map2 (⊔) As Bs) of
        Some xs ⇒ Some (Prule α xs)
      | None ⇒ None)
    else None)
| Prule α As ⊔ B =
    (case match B (to-pterm (lhs α)) of
      None ⇒ None
    | Some σ ⇒
      (case those (map2 (⊔) As (map σ (var-rule α))) of
        Some xs ⇒ Some (Prule α xs)
      | None ⇒ None))
| A ⊔ Prule α Bs =
    (case match A (to-pterm (lhs α)) of
      None ⇒ None
    | Some σ ⇒
      (case those (map2 (⊔) (map σ (var-rule α)) Bs) of
        Some xs ⇒ Some (Prule α xs)
      | None ⇒ None))
| A ⊔ B = None

lemma join-sym: A ⊔ B = B ⊔ A
proof(induct rule:join.induct)
  case (2 f As g Bs)
  then show ?case proof(cases f = g ∧ length As = length Bs)
    case True
    with 2 have ∀(a,b) ∈ set (zip As Bs). a ⊔ b = b ⊔ a
      by auto
    with True have (map2 (⊔) As Bs) = (map2 (⊔) Bs As)
      by (smt case-prod-unfold map-eq-conv' map-fst-zip map-snd-zip nth-mem)
    then show ?thesis using 2 unfolding join.simps
      by auto
  qed auto
next
  case (3 α As β Bs)
  then show ?case proof(cases α = β)
    case True
```

with *3* **have** ∗:∀ (*a*,*b*) ∈ *set* (*zip As Bs*). *a* ⊔ *b* = *b* ⊔ *a*
by *auto*
have *length* (*map2* (⊔) *As Bs*) = *length* (*map2* (⊔) *Bs As*)
by *auto*
with ∗ **have** (*map2* (⊔) *As Bs*) = (*map2* (⊔) *Bs As*)
by (*smt fst-conv length-map length-zip map-eq-conv′ min-less-iff-conj nth-mem nth-zip prod.case-eq-if snd-conv*)
then show *?thesis* **using** *3* **unfolding** *join.simps*
by *auto*
qed *auto*
next
case (*4-1 α As v*)
then show *?case* **proof**(*cases match* (*Var v*) (*to-pterm* (*lhs α*)) = *None*)
case *False*
then obtain *σ* **where** *sigma*:*match* (*Var v*) (*to-pterm* (*lhs α*)) = *Some σ*
by *blast*
with *4-1* **have** ∗:∀ (*a*,*b*) ∈ *set* (*zip As* (*map σ* (*var-rule α*))). *a* ⊔ *b* = *b* ⊔ *a*
by *auto*
have *length* (*map2* (⊔) *As* (*map σ* (*var-rule α*))) = *length* (*map2* (⊔) (*map σ* (*var-rule α*)) *As*)
by *auto*
with ∗ **have** (*map2* (⊔) *As* (*map σ* (*var-rule α*))) = (*map2* (⊔) (*map σ* (*var-rule α*)) *As*)
by (*smt fst-conv length-map length-zip map-eq-conv′ min-less-iff-conj nth-mem nth-zip prod.case-eq-if snd-conv*)
then show *?thesis* **unfolding** *join.simps sigma*
by *simp*
qed *simp*
next
case (*4-2 α As f Bs*)
then show *?case* **proof**(*cases match* (*Pfun f Bs*) (*to-pterm* (*lhs α*)) = *None*)
case *False*
then obtain *σ* **where** *sigma*:*match* (*Pfun f Bs*) (*to-pterm* (*lhs α*)) = *Some σ*
by *blast*
with *4-2* **have** ∗:∀ (*a*,*b*) ∈ *set* (*zip As* (*map σ* (*var-rule α*))). *a* ⊔ *b* = *b* ⊔ *a*
by *auto*
have *length* (*map2* (⊔) *As* (*map σ* (*var-rule α*))) = *length* (*map2* (⊔) (*map σ* (*var-rule α*)) *As*)
by *auto*
with ∗ **have** (*map2* (⊔) *As* (*map σ* (*var-rule α*))) = (*map2* (⊔) (*map σ* (*var-rule α*)) *As*)
by (*smt fst-conv length-map length-zip map-eq-conv′ min-less-iff-conj nth-mem nth-zip prod.case-eq-if snd-conv*)
then show *?thesis* **unfolding** *join.simps sigma*
by *simp*
qed *simp*
next
case (*5-1 v α Bs*)
then show *?case* **proof**(*cases match* (*Var v*) (*to-pterm* (*lhs α*)) = *None*)

case *False*
then obtain σ **where** *sigma*:*match* (*Var v*) (*to-pterm* (*lhs* α)) = *Some* σ
by *blast*
with *5-1* **have** $*$:$\forall$(*a*,*b*) $\in$ *set* (*zip* (*map* σ (*var-rule* α)) *Bs*). *a* $\sqcup$ *b* = *b* $\sqcup$ *a*
by *auto*
have *length* (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *Bs*) = *length* (*map2* ($\sqcup$) *Bs* (*map* σ (*var-rule* α)))
by *auto*
with $*$ **have** (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *Bs*) = (*map2* ($\sqcup$) *Bs* (*map* σ (*var-rule* α)))
by (*smt fst-conv length-map length-zip map-eq-conv′ min-less-iff-conj nth-mem nth-zip prod.case-eq-if snd-conv*)
then show *?thesis* **unfolding** *join.simps sigma*
by *simp*
qed *simp*
next
case (*5-2 f As* α *Bs*)
then show *?case* **proof**(*cases match* (*Pfun f As*) (*to-pterm* (*lhs* α)) = *None*)
case *False*
then obtain σ **where** *sigma*:*match* (*Pfun f As*) (*to-pterm* (*lhs* α)) = *Some* σ
by *blast*
with *5-2* **have** $*$:$\forall$(*a*,*b*) $\in$ *set* (*zip* (*map* σ (*var-rule* α)) *Bs*). *a* $\sqcup$ *b* = *b* $\sqcup$ *a*
by *auto*
have *length* (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *Bs*) = *length* (*map2* ($\sqcup$) *Bs* (*map* σ (*var-rule* α)))
by *auto*
with $*$ **have** (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *Bs*) = (*map2* ($\sqcup$) *Bs* (*map* σ (*var-rule* α)))
by (*smt fst-conv length-map length-zip map-eq-conv′ min-less-iff-conj nth-mem nth-zip prod.case-eq-if snd-conv*)
then show *?thesis* **unfolding** *join.simps sigma*
by *simp*
qed *simp*
qed *simp-all*

lemma *join-with-source*:
assumes *A* $\in$ *wf-pterm R*
shows *A* $\sqcup$ *to-pterm* (*source A*) = *Some A*
using *assms* **proof**(*induct A*)
case (*2 As f*)
then have $\forall i <$ *length As*. (*map2* ($\sqcup$) *As* (*map to-pterm* (*map source As*)))!*i* = *Some* (*As*!*i*)
by *simp*
then have *those* (*map2* ($\sqcup$) *As* (*map to-pterm* (*map source As*))) = *Some As*
by (*simp add*: *those-some*)
then show *?case*
by *simp*
next
case (*3* α *As*)

then have $\forall\, i < \mathit{length\ As}.\ (\mathit{map2}\ (\sqcup)\ \mathit{As}\ (\mathit{map\ to\text{-}pterm}\ (\mathit{map\ source\ As})))!i = \mathit{Some}\ (\mathit{As}!i)$
by *simp*
then have $\mathit{IH}{:}\mathit{those}\ (\mathit{map2}\ (\sqcup)\ \mathit{As}\ (\mathit{map\ to\text{-}pterm}\ (\mathit{map\ source\ As}))) = \mathit{Some\ As}$
by (*simp add*: *those-some*)
from *3*(*1*) **have** $\mathit{match}{:}\mathit{match}\ (\mathit{to\text{-}pterm}\ (\mathit{source}\ (\mathit{Prule}\ \alpha\ \mathit{As})))\ (\mathit{to\text{-}pterm}\ (\mathit{lhs}\ \alpha)) = \mathit{Some}\ (\langle \mathit{map}\ (\mathit{to\text{-}pterm} \circ \mathit{source})\ \mathit{As}\rangle_\alpha)$
by (*metis* (*no-types*, *lifting*) *fun-mk-subst lhs-subst-trivial map-map source.simps*(*3*) *to-pterm.simps*(*1*) *to-pterm-subst*)
from *3*(*2*) **have** $(\mathit{map}\ (\langle \mathit{map}\ (\mathit{to\text{-}pterm} \circ \mathit{source})\ \mathit{As}\rangle_\alpha)\ (\mathit{var\text{-}rule}\ \alpha)) = \mathit{map}\ (\mathit{to\text{-}pterm} \circ \mathit{source})\ \mathit{As}$
by (*metis apply-lhs-subst-var-rule length-map*)
with *IH match join.simps*(*4*,*5*) **show** *?case* **by**(*cases source* (*Prule* α *As*)) *simp-all*
qed *simp*

context *no-var-lhs*
begin

lemma *join-subst*:
assumes $B \in \mathit{wf\text{-}pterm}\ R$ **and** $\forall\, x \in \mathit{vars\text{-}term}\ B.\ \varrho\ x \in \mathit{wf\text{-}pterm}\ R$
and $\forall\, x \in \mathit{vars\text{-}term}\ B.\ \mathit{source}\ (\varrho\ x) = \sigma\ x$
shows $(B \cdot (\mathit{to\text{-}pterm} \circ \sigma)) \sqcup ((\mathit{to\text{-}pterm}\ (\mathit{source}\ B)) \cdot \varrho) = \mathit{Some}\ (B \cdot \varrho)$
using *assms* **proof**(*induct B*)
case (*1 x*)
then show *?case* **unfolding** *eval-term.simps source.simps to-pterm.simps o-apply*
using *join-with-source* **by** (*metis term.set-intros*(*3*) *join-sym*)
next
case (*2 ts f*)
{fix *i* **assume** $i{:}i < \mathit{length\ ts}$
with *2* **have** $((\mathit{ts}!i) \cdot (\mathit{to\text{-}pterm} \circ \sigma)) \sqcup ((\mathit{to\text{-}pterm}\ (\mathit{source}\ (\mathit{ts}!i))) \cdot \varrho) = \mathit{Some}\ (\mathit{ts}!i \cdot \varrho)$
by (*meson nth-mem term.set-intros*(*4*))
then have $\mathit{map2}\ (\sqcup)\ (\mathit{map}\ (\lambda t.\ t \cdot (\mathit{to\text{-}pterm} \circ \sigma))\ \mathit{ts})\ (\mathit{map}\ (\lambda t.\ t \cdot \varrho)\ (\mathit{map\ to\text{-}pterm}\ (\mathit{map\ source\ ts})))!i = \mathit{Some}\ ((\mathit{map}\ (\lambda t.\ t \cdot \varrho)\ \mathit{ts})!i)$
using *i* **by** *fastforce*
}
then have $\mathit{those}\ (\mathit{map2}\ (\sqcup)\ (\mathit{map}\ (\lambda t.\ t \cdot (\mathit{to\text{-}pterm} \circ \sigma))\ \mathit{ts})\ (\mathit{map}\ (\lambda t.\ t \cdot \varrho)\ (\mathit{map\ to\text{-}pterm}\ (\mathit{map\ source\ ts})))) = \mathit{Some}\ (\mathit{map}\ (\lambda t.\ t \cdot \varrho)\ \mathit{ts})$
using *those-some* **by** (*smt* (*verit*) *length-map length-zip min.idem*)
then show *?case*
unfolding *source.simps to-pterm.simps eval-term.simps* **using** *join.simps*(*2*) **by** *auto*
next
case (*3* α *As*)
from *3*(*1*) *no-var-lhs* **obtain** *f ts* **where** $f{:}\mathit{lhs}\ \alpha = \mathit{Fun}\ f\ \mathit{ts}$
by *fastforce*
obtain τ **where** $\mathit{match}{:}\mathit{match}\ (\mathit{to\text{-}pterm}\ (\mathit{lhs}\ \alpha \cdot \langle \mathit{map\ source\ As}\rangle_\alpha) \cdot \varrho)\ (\mathit{to\text{-}pterm}\ (\mathit{lhs}\ \alpha)) = \mathit{Some}\ \tau$

and τ:($\forall\, x{\in}$*vars-term* (*lhs* α). $\tau\ x$ = ((*to-pterm* $\circ$ $\langle$*map source As*$\rangle_\alpha$) $\circ_s$ ϱ) x)
using *match-complete′* **unfolding** *to-pterm-subst* **by** (*smt* (*verit*, *best*) *set-vars-term-list subst-subst vars-to-pterm*)
{fix *i* **assume** *i*:i < *length* (*var-rule* α)
let *?x* =*var-rule* α ! *i*
have ((*to-pterm* $\circ$ $\langle$*map source As*$\rangle_\alpha$) $\circ_s$ ϱ) *?x* = *to-pterm* (*source* (*As*!*i*)) $\cdot$ ϱ
using *i* *3*(*2*) **by** (*simp add*: *mk-subst-distinct subst-compose-def*)
moreover from *3* **have** ((*As*!*i*) $\cdot$ (*to-pterm* $\circ$ σ)) $\sqcup$ (*to-pterm* (*source* (*As*!*i*)) $\cdot$ ϱ) = *Some* ((*As*!*i*) $\cdot$ ϱ)
by (*metis* (*mono-tags*, *lifting*) *i nth-mem term.set-intros*(*4*))
ultimately have ((*As*!*i*) $\cdot$ (*to-pterm* $\circ$ σ)) $\sqcup$ (τ *?x*) = *Some* ((*As*!*i*) $\cdot$ ϱ)
using τ **by** (*metis comp-apply i nth-mem set-remdups set-rev set-vars-term-list*)
then have (*map2* ($\sqcup$) (*map* ($\lambda t.\ t \cdot$ (*to-pterm* $\circ$ σ)) *As*) (*map* τ (*var-rule* α)))!*i* = *Some* ((*map* ($\lambda t.\ t \cdot \varrho$) *As*)!*i*)
using *3*(*2*) *i* **by** *auto*
}
then have *those* (*map2* ($\sqcup$) (*map* ($\lambda t.\ t \cdot$ (*to-pterm* $\circ$ σ)) *As*) (*map* τ (*var-rule* α))) = *Some* (*map* ($\lambda t.\ t \cdot \varrho$) *As*)
by (*simp add*: *3*(*2*) *those-some*)
then show *?case*
using *match* **unfolding** *source.simps to-pterm.simps eval-term.simps f* **using** *join.simps*(*5*) *f* **by** *auto*
qed

end

lemma *join-same*:
shows $A \sqcup A$ = *Some* A
proof(*induct A*)
case (*Pfun f As*)
{fix *i* **assume** *i*:i < *length As*
with *Pfun* **have** *As*!*i* $\sqcup$ *As*!*i* = *Some* (*As*!*i*) **by** *simp*
with *i* **have** (*map2* ($\sqcup$) *As As*)!*i* = *Some* (*As*!*i*) **by** *simp*
}
then have *those* (*map2* ($\sqcup$) *As As*) = *Some As*
by (*simp add*: *those-some*)
then show *?case* **unfolding** *join.simps* **by** *simp*
next
case (*Prule* α *As*)
{fix *i* **assume** *i*:i < *length As*
with *Prule* **have** *As*!*i* $\sqcup$ *As*!*i* = *Some* (*As*!*i*) **by** *simp*
with *i* **have** (*map2* ($\sqcup$) *As As*)!*i* = *Some* (*As*!*i*) **by** *simp*
}
then have *those* (*map2* ($\sqcup$) *As As*) = *Some As*
by (*simp add*: *those-some*)
then show *?case* **unfolding** *join.simps* **by** *simp*
qed *simp*

Analogous to residuals there are 6 lemmas corresponding to the step cases in induction proofs for joins.

lemma *join-fun-fun*:
  assumes $(\mathit{Pfun}\ f\ \mathit{As}) \sqcup (\mathit{Pfun}\ g\ \mathit{Bs}) = \mathit{Some}\ C$
  shows $f = g \wedge \mathit{length}\ \mathit{As} = \mathit{length}\ \mathit{Bs} \wedge$
    $(\exists\, \mathit{Cs}.\ C = \mathit{Pfun}\ f\ \mathit{Cs} \wedge$
    $\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As} \wedge$
    $(\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup \mathit{Bs}!i = \mathit{Some}\ (\mathit{Cs}!i)))$
proof−
  have $*$:$f = g \wedge \mathit{length}\ \mathit{As} = \mathit{length}\ \mathit{Bs}$
    using *assms join.simps*(*2*) **by** (*metis option.simps*(*3*))
  then obtain *Cs* **where** *Cs*:$\mathit{those}\ (\mathit{map2}\ (\sqcup)\ \mathit{As}\ \mathit{Bs}) = \mathit{Some}\ \mathit{Cs}$
    using *assms option.simps*(*3*) *option.simps*(*4*) **by** *fastforce*
  hence $\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup \mathit{Bs}!i = \mathit{Some}\ (\mathit{Cs}!i)$
    using $*$ *those-some2* **by** *fastforce*
  with $*$ *Cs assms*(*1*) **show** *?thesis*
    using *length-those* **by** *fastforce*
qed

lemma *join-rule-rule*:
  assumes $(\mathit{Prule}\ \alpha\ \mathit{As}) \sqcup (\mathit{Prule}\ \beta\ \mathit{Bs}) = \mathit{Some}\ C$
    $(\mathit{Prule}\ \alpha\ \mathit{As}) \in \mathit{wf\text{-}pterm}\ R$
    $(\mathit{Prule}\ \beta\ \mathit{Bs}) \in \mathit{wf\text{-}pterm}\ R$
  shows $\alpha = \beta \wedge \mathit{length}\ \mathit{As} = \mathit{length}\ \mathit{Bs} \wedge$
    $(\exists\, \mathit{Cs}.\ C = \mathit{Prule}\ \alpha\ \mathit{Cs} \wedge$
    $\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As} \wedge$
    $(\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup \mathit{Bs}!i = \mathit{Some}\ (\mathit{Cs}!i)))$
proof−
  have $\alpha = \beta$
    using *assms*(*1*) *join.simps*(*3*) **by** (*metis option.simps*(*3*))
  with *assms*(*2*,*3*) **have** *l*: $\mathit{length}\ \mathit{As} = \mathit{length}\ \mathit{Bs}$
    using *length-args-well-Prule* **by** *blast*
  from ‹$\alpha = \beta$› **obtain** *Cs* **where** *Cs*:$\mathit{those}\ (\mathit{map2}\ (\sqcup)\ \mathit{As}\ \mathit{Bs}) = \mathit{Some}\ \mathit{Cs}$
    using *assms* **by** *fastforce*
  hence $\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup \mathit{Bs}!i = \mathit{Some}\ (\mathit{Cs}!i)$
    using *l those-some2* **by** *fastforce*
  with ‹$\alpha = \beta$› *l Cs assms*(*1*) **show** *?thesis*
    using *length-those* **by** *fastforce*
qed

lemma *join-rule-var*:
  assumes $(\mathit{Prule}\ \alpha\ \mathit{As}) \sqcup (\mathit{Var}\ x) = \mathit{Some}\ C$
    $(\mathit{Prule}\ \alpha\ \mathit{As}) \in \mathit{wf\text{-}pterm}\ R$
  shows $\exists \sigma.\ \mathit{match}\ (\mathit{Var}\ x)\ (\mathit{to\text{-}pterm}\ (\mathit{lhs}\ \alpha)) = \mathit{Some}\ \sigma \wedge$
    $(\exists\, \mathit{Cs}.\ C = \mathit{Prule}\ \alpha\ \mathit{Cs} \wedge$
    $\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As} \wedge$
    $(\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup (\sigma\ (\mathit{var\text{-}rule}\ \alpha\ !\ i)) = \mathit{Some}\ (\mathit{Cs}!i)))$
proof−
  from *assms*(*2*) **have** *l*:$\mathit{length}\ \mathit{As} = \mathit{length}\ (\mathit{var\text{-}rule}\ \alpha)$
    using *wf-pterm.cases* **by** *auto*

obtain *σ* **where** *σ*:*match* (*Var x*) (*to-pterm* (*lhs α*)) = *Some σ*
using *assms*(*1*) **by** *fastforce*
then obtain *Cs* **where** *Cs*:*those* (*map2* (⊔) *As* (*map σ* (*var-rule α*))) = *Some Cs*
using *assms*(*1*) **by** *fastforce*
with *l* **have** *l2*:*length Cs* = *length As*
using *length-those* **by** *fastforce*
from *Cs* **have** ∀ *i* < *length As*. *As*!*i* ⊔ (*σ* (*var-rule α* ! *i*)) = *Some* (*Cs*!*i*)
using *l those-some2* **by** *fastforce*
with *σ Cs assms*(*1*) *l2* **show** *?thesis* **by** *simp*
qed

lemma *join-rule-fun*:
assumes (*Prule α As*) ⊔ (*Pfun f Bs*) = *Some C*
(*Prule α As*) ∈ *wf-pterm R*
shows ∃ *σ*. *match* (*Pfun f Bs*) (*to-pterm* (*lhs α*)) = *Some σ* ∧
(∃ *Cs*. *C* = *Prule α Cs* ∧
length Cs = *length As* ∧
(∀ *i* < *length As*. *As*!*i* ⊔ (*σ* (*var-rule α* ! *i*)) = *Some* (*Cs*!*i*)))
proof−
from *assms*(*2*) **have** *l*:*length As* = *length* (*var-rule α*)
using *wf-pterm.simps* **by** *fastforce*
obtain *σ* **where** *σ*:*match* (*Pfun f Bs*) (*to-pterm* (*lhs α*)) = *Some σ*
using *assms*(*1*) **by** *fastforce*
then obtain *Cs* **where** *Cs*:*those* (*map2* (⊔) *As* (*map σ* (*var-rule α*))) = *Some Cs*
using *assms*(*1*) **by** *fastforce*
with *l* **have** *l2*:*length Cs* = *length As*
using *length-those* **by** *fastforce*
from *Cs* **have** ∀ *i* < *length As*. *As*!*i* ⊔ (*σ* (*var-rule α* ! *i*)) = *Some* (*Cs*!*i*)
using *l those-some2* **by** *fastforce*
with *σ Cs assms*(*1*) *l2* **show** *?thesis* **by** *auto*
qed

lemma *join-wf-pterm*:
assumes *A* ⊔ *B* = *Some C*
and *A* ∈ *wf-pterm R* **and** *B* ∈ *wf-pterm R*
shows *C* ∈ *wf-pterm R*
using *assms* **proof**(*induct A B arbitrary*:*C rule*:*join.induct*)
case (*1 x y*)
then show *?case*
by (*metis join.simps*(*1*) *option.distinct*(*1*) *option.sel*)
next
case (*2 f As g Bs*)
then obtain *Cs* **where** *f* = *g* ∧ *length As* = *length Bs* ∧
C = *Pfun f Cs* ∧
length Cs = *length As* ∧
(∀ *i* < *length As*. *As*!*i* ⊔ *Bs*!*i* = *Some* (*Cs*!*i*))
by (*meson join-fun-fun*)

moreover with *2* **have** $i < \mathit{length}\ \mathit{As} \Longrightarrow \mathit{Cs}!i \in \mathit{wf\text{-}pterm}\ R$ **for** i
using *fun-well-arg* **by** (*metis* (*no-types*, *opaque-lifting*) *fst-conv in-set-zip nth-mem snd-conv*)
ultimately show *?case*
by (*metis in-set-conv-nth wf-pterm.intros*(*2*))
next
case ($3\ \alpha\ \mathit{As}\ \beta\ \mathit{Bs}$)
then obtain *Cs* **where** $\alpha = \beta \wedge \mathit{length}\ \mathit{As} = \mathit{length}\ \mathit{Bs} \wedge$
$(C = \mathit{Prule}\ \alpha\ \mathit{Cs} \wedge$
$\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As} \wedge$
$(\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup \mathit{Bs}!i = \mathit{Some}\ (\mathit{Cs}!i)))$
by (*meson join-rule-rule*)
moreover with *3* **have** $i < \mathit{length}\ \mathit{As} \Longrightarrow \mathit{Cs}!i \in \mathit{wf\text{-}pterm}\ R$ **for** i
using *fun-well-arg* **by** (*metis* (*no-types*, *opaque-lifting*) *fst-conv in-set-zip nth-mem snd-conv*)
moreover from *3*(*3*) **have** $\mathit{to\text{-}rule}\ \alpha \in R$
using *wf-pterm.simps* **by** *fastforce*
ultimately show *?case*
by (*smt* (*verit*, *best*) *3.prems*(*2*) *in-set-idx old.sum.distinct*(*2*) *term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.cases wf-pterm.intros*(*3*))
next
case (*4-1* α *As* v)
then obtain $\mathit{Cs}\ \sigma$ **where** $*{:}\mathit{match}\ (\mathit{Var}\ v)\ (\mathit{to\text{-}pterm}\ (\mathit{lhs}\ \alpha)) = \mathit{Some}\ \sigma \wedge$
$C = \mathit{Prule}\ \alpha\ \mathit{Cs} \wedge$
$\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As} \wedge$
$(\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup (\sigma\ (\mathit{var\text{-}rule}\ \alpha\ !\ i)) = \mathit{Some}\ (\mathit{Cs}!i))$
by (*meson join-rule-var*)
from *4-1*(*3*) **have** $\mathit{wellA}{:}\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \in \mathit{wf\text{-}pterm}\ R$
by *auto*
from *4-1*(*3*) **have** l: $\mathit{length}\ \mathit{As} = \mathit{length}\ (\mathit{var\text{-}rule}\ \alpha)$
using *wf-pterm.simps* **by** *fastforce*
then have $\mathit{l2}{:}\mathit{length}\ \mathit{As} = \mathit{length}\ (\mathit{zip}\ \mathit{As}\ (\mathit{map}\ \sigma\ (\mathit{var\text{-}rule}\ \alpha)))$
by *simp*
from $l\ *$ **have** $\forall\, i < \mathit{length}\ \mathit{As}.\ \sigma\ (\mathit{var\text{-}rule}\ \alpha\ !\ i) \in \mathit{wf\text{-}pterm}\ R$
using *4-1*(*4*) **by** (*metis match-well-def vars-to-pterm*)
with *4-1*(*1*) $*$ *wellA l2* **have** $\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{Cs}!i \in \mathit{wf\text{-}pterm}\ R$
by (*smt* (*z3*) *l length-map nth-map nth-mem nth-zip*)
with *4-1*(*3*) $*$ **show** *?case*
by (*smt* (*verit*) *Inr-not-Inl in-set-conv-nth term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.cases wf-pterm.intros*(*3*))
next
case (*4-2* α *As f Bs*)
then obtain $\sigma\ \mathit{Cs}$ **where** $*{:}\mathit{match}\ (\mathit{Pfun}\ f\ \mathit{Bs})\ (\mathit{to\text{-}pterm}\ (\mathit{lhs}\ \alpha)) = \mathit{Some}\ \sigma \wedge$
$(C = \mathit{Prule}\ \alpha\ \mathit{Cs} \wedge$
$\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As} \wedge$
$(\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup (\sigma\ (\mathit{var\text{-}rule}\ \alpha\ !\ i)) = \mathit{Some}\ (\mathit{Cs}!i)))$
by (*meson join-rule-fun*)
from *4-2*(*3*) **have** $\mathit{wellA}{:}\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \in \mathit{wf\text{-}pterm}\ R$
by *auto*

from *4-2(3)* **have** *l*: *length As* = *length* (*var-rule* α)
  using *wf-pterm.simps* **by** *fastforce*
then have *l2*:*length As* = *length* (*zip As* (*map* σ (*var-rule* α)))
  by *simp*
from *l* $*$ **have** $\forall$ *i* < *length As*. σ (*var-rule* α ! *i*) $\in$ *wf-pterm R*
  using *4-2(4)* **by** (*metis match-well-def vars-to-pterm*)
with *4-2(1)* $*$ *wellA l2* **have** $\forall$ *i* < *length As*. *Cs*!*i* $\in$ *wf-pterm R*
  by (*smt l length-map nth-map nth-mem nth-zip*)
with *4-2(3)* $*$ **show** *?case*
   by (*smt* (*verit*, *ccfv-threshold*) *Inr-not-Inl in-set-conv-nth term.distinct(1) term.inject(2) wf-pterm.cases wf-pterm.intros(3)*)
next
case (*5-1 v* α *As*)
then obtain *Cs* σ **where** $*$:*match* (*Var v*) (*to-pterm* (*lhs* α)) = *Some* σ $\wedge$
  C = *Prule* α *Cs* $\wedge$
  length Cs = *length As* $\wedge$
  ($\forall$ *i* < *length As*. (σ (*var-rule* α ! *i*)) $\sqcup$ *As*!*i* = *Some* (*Cs*!*i*))
  using *join-rule-var* **by** (*metis join-sym*)
from *5-1(4)* **have** *wellA*:$\forall$ *i* < *length As*. *As*!*i* $\in$ *wf-pterm R*
  by *auto*
from *5-1(4)* **have** *l*: *length As* = *length* (*var-rule* α)
  using *wf-pterm.simps* **by** *fastforce*
then have *l2*:*length As* = *length* (*zip As* (*map* σ (*var-rule* α)))
  by *simp*
from *l* $*$ **have** $\forall$ *i* < *length As*. σ (*var-rule* α ! *i*) $\in$ *wf-pterm R*
  using *5-1(3)* **by** (*metis match-well-def vars-to-pterm*)
with *5-1(1)* $*$ *wellA l2 l* **have** $\forall$ *i* < *length As*. *Cs*!*i* $\in$ *wf-pterm R*
  by (*smt* (*verit*, *del-insts*) *length-map nth-map nth-mem nth-zip zip-symm*)
with *5-1(4)* $*$ **show** *?case*
   by (*smt* (*verit*) *Inr-not-Inl in-set-conv-nth term.distinct(1) term.inject(2) wf-pterm.cases wf-pterm.intros(3)*)
next
case (*5-2 f Bs* α *As*)
then obtain *Cs* σ **where** $*$:*match* (*Pfun f Bs*) (*to-pterm* (*lhs* α)) = *Some* σ $\wedge$
  C = *Prule* α *Cs* $\wedge$
  length Cs = *length As* $\wedge$
  ($\forall$ *i* < *length As*. (σ (*var-rule* α ! *i*)) $\sqcup$ *As*!*i* = *Some* (*Cs*!*i*))
  using *join-sym join-rule-fun* **by** *metis*
from *5-2(4)* **have** *wellA*:$\forall$ *i* < *length As*. *As*!*i* $\in$ *wf-pterm R*
  by *auto*
from *5-2(4)* **have** *l*: *length As* = *length* (*var-rule* α)
  using *wf-pterm.simps* **by** *fastforce*
then have *l2*:*length As* = *length* (*zip* (*map* σ (*var-rule* α)) *As*)
  by *simp*
from *l* $*$ **have** $\forall$ *i* < *length As*. σ (*var-rule* α ! *i*) $\in$ *wf-pterm R*
  using *5-2(3)* **by** (*metis match-well-def vars-to-pterm*)
with *5-2(1)* $*$ *wellA l2* **have** $\forall$ *i* < *length As*. *Cs*!*i* $\in$ *wf-pterm R*
  by (*smt l length-map nth-map nth-mem nth-zip*)
with $*$ **show** *?case*

by (*metis 5-2.prems*(*3*) *Inl-inject Inr-not-Inl in-set-idx l term.distinct*(*1*) *term.sel*(*2*) *wf-pterm.cases wf-pterm.intros*(*3*))
qed *auto*

lemma *source-join*:
assumes $A \sqcup B = \mathit{Some}\ C$
and $A \in \mathit{wf\text{-}pterm}\ R$ **and** $B \in \mathit{wf\text{-}pterm}\ R$
shows *co-initial* $A\ C$
using *assms* **proof**(*induct* $A\ B$ *arbitrary*:C *rule*:*join.induct*)
case (*1 x y*)
then show *?case*
by (*metis join.simps*(*1*) *option.discI option.sel*)
next
case (*2 f As g Bs*)
then obtain *Cs* **where** $f = g \wedge \mathit{length}\ \mathit{As} = \mathit{length}\ \mathit{Bs} \wedge$
$C = \mathit{Pfun}\ f\ \mathit{Cs} \wedge$
$\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As} \wedge$
$(\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup \mathit{Bs}!i = \mathit{Some}\ (\mathit{Cs}!i))$
by (*meson join-fun-fun*)
moreover with *2* **have** $i < \mathit{length}\ \mathit{As} \Longrightarrow$ *co-initial* $(\mathit{As}!i)\ (\mathit{Cs}!i)$ **for** i
using *fun-well-arg* **by** (*metis* (*no-types*, *opaque-lifting*) *fst-conv in-set-zip nth-mem snd-conv*)
ultimately show *?case*
by (*simp add*: *nth-map-conv*)
next
case (*3* α *As* β *Bs*)
then obtain *Cs* **where** $\alpha = \beta \wedge \mathit{length}\ \mathit{As} = \mathit{length}\ \mathit{Bs} \wedge$
$(C = \mathit{Prule}\ \alpha\ \mathit{Cs} \wedge$
$\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As} \wedge$
$(\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup \mathit{Bs}!i = \mathit{Some}\ (\mathit{Cs}!i)))$
by (*meson join-rule-rule*)
moreover with *3* **have** $i < \mathit{length}\ \mathit{As} \Longrightarrow$ *co-initial* $(\mathit{As}!i)\ (\mathit{Cs}!i)$ **for** i
using *fun-well-arg* **by** (*metis* (*no-types*, *opaque-lifting*) *fst-conv in-set-zip nth-mem snd-conv*)
ultimately show *?case*
by (*metis* (*mono-tags*, *lifting*) *map-eq-conv′ source.simps*(*3*))
next
case (*4-1* α *As v*)
then obtain *Cs* σ **where** $*$:*match* (*Var v*) (*to-pterm* (*lhs* α)) $= \mathit{Some}\ \sigma \wedge$
$C = \mathit{Prule}\ \alpha\ \mathit{Cs} \wedge$
$\mathit{length}\ \mathit{Cs} = \mathit{length}\ \mathit{As} \wedge$
$(\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \sqcup (\sigma\ (\mathit{var\text{-}rule}\ \alpha\ !\ i)) = \mathit{Some}\ (\mathit{Cs}!i))$
by (*meson join-rule-var*)
from *4-1*(*3*) **have** *wellA*:$\forall\, i < \mathit{length}\ \mathit{As}.\ \mathit{As}!i \in \mathit{wf\text{-}pterm}\ R$
by *auto*
from *4-1*(*3*) **have** *l*: $\mathit{length}\ \mathit{As} = \mathit{length}\ (\mathit{var\text{-}rule}\ \alpha)$
using *wf-pterm.simps* **by** *fastforce*
then have *l2*:$\mathit{length}\ \mathit{As} = \mathit{length}\ (\mathit{zip}\ \mathit{As}\ (\mathit{map}\ \sigma\ (\mathit{var\text{-}rule}\ \alpha)))$
by *simp*

from l $*$ **have** $\forall\, i < length\ As.\ \sigma\ (var\text{-}rule\ \alpha\ !\ i) \in wf\text{-}pterm\ R$
using *4-1*(*4*) **by** (*metis match-well-def vars-to-pterm*)
with *4-1*(*1*) $*$ *wellA l2* **have** $\forall\, i < length\ As.\ co\text{-}initial\ (As!i)\ (Cs!i)$
by (*smt* (*z3*) *l length-map nth-map nth-mem nth-zip*)
with *4-1*(*3*) $*$ **show** *?case*
by (*metis nth-map-conv source.simps*(*3*))
next
case (*4-2* α *As f Bs*)
then obtain σ *Cs* **where** $*{:}match\ (Pfun\ f\ Bs)\ (to\text{-}pterm\ (lhs\ \alpha)) = Some\ \sigma\ \wedge$
$(C = Prule\ \alpha\ Cs\ \wedge$
$length\ Cs = length\ As\ \wedge$
$(\forall\, i < length\ As.\ As!i \sqcup (\sigma\ (var\text{-}rule\ \alpha\ !\ i)) = Some\ (Cs!i)))$
by (*meson join-rule-fun*)
from *4-2*(*3*) **have** $wellA{:}\forall\, i < length\ As.\ As!i \in wf\text{-}pterm\ R$
by *auto*
from *4-2*(*3*) **have** l: $length\ As = length\ (var\text{-}rule\ \alpha)$
using *wf-pterm.simps* **by** *fastforce*
then have $l2{:}length\ As = length\ (zip\ As\ (map\ \sigma\ (var\text{-}rule\ \alpha)))$
by *simp*
from l $*$ **have** $\forall\, i < length\ As.\ \sigma\ (var\text{-}rule\ \alpha\ !\ i) \in wf\text{-}pterm\ R$
using *4-2*(*4*) **by** (*metis match-well-def vars-to-pterm*)
with *4-2*(*1*) $*$ *wellA l2* **have** $\forall\, i < length\ As.\ co\text{-}initial\ (As!i)\ (Cs!i)$
by (*smt l length-map nth-map nth-mem nth-zip*)
with *4-2*(*3*) $*$ **show** *?case*
by (*metis nth-map-conv source.simps*(*3*))
next
case (*5-1 v* α *As*)
then obtain *Cs* σ **where** $*{:}match\ (Var\ v)\ (to\text{-}pterm\ (lhs\ \alpha)) = Some\ \sigma$
$C = Prule\ \alpha\ Cs$
$length\ Cs = length\ As\ \wedge$
$(\forall\, i < length\ As.\ (\sigma\ (var\text{-}rule\ \alpha\ !\ i)) \sqcup As!i = Some\ (Cs!i))$
using *join-rule-var* **by** (*metis join-sym*)
from *5-1*(*4*) **have** $wellA{:}\forall\, i < length\ As.\ As!i \in wf\text{-}pterm\ R$
by *auto*
from *5-1*(*4*) **have** l: $length\ As = length\ (var\text{-}rule\ \alpha)$
using *wf-pterm.simps* **by** *fastforce*
then have $l2{:}length\ As = length\ (zip\ As\ (map\ \sigma\ (var\text{-}rule\ \alpha)))$
by *simp*
from l $*$ **have** $\forall\, i < length\ As.\ \sigma\ (var\text{-}rule\ \alpha\ !\ i) \in wf\text{-}pterm\ R$
using *5-1*(*3*) **by** (*metis match-well-def vars-to-pterm*)
with *5-1*(*1*) $*$ *wellA l2 l* **have** $IH{:}\forall\, i < length\ As.\ co\text{-}initial\ ((map\ \sigma\ (var\text{-}rule\ \alpha))!i)\ (Cs!i)$
by (*smt* (*verit*, *del-insts*) *length-map nth-map nth-mem nth-zip zip-symm*)
moreover have $v{:}\,Var\ v = (to\text{-}pterm\ (lhs\ \alpha)) \cdot \langle(map\ \sigma\ (var\text{-}rule\ \alpha))\rangle_\alpha$
using $*$ **by** (*smt* (*verit*, *ccfv-threshold*) *apply-lhs-subst-var-rule map-eq-conv match-lhs-subst*)
show *?case* **using** *IH l* **unfolding** v $*$(*2*) *source.simps*
by (*metis* $*$(*1*,*3*) *fun-mk-subst length-map lhs-subst-trivial nth-map-conv option.inject source.simps*(*1*) *source-apply-subst source-to-pterm to-pterm-wf-pterm*

v)
next
case (*5-2 f Bs* α *As*)
then obtain *Cs* σ **where** $*$:*match* (*Pfun f Bs*) (*to-pterm* (*lhs* α)) = *Some* σ
C = *Prule* α *Cs*
length Cs = *length As* $\wedge$
($\forall\, i$ < *length As*. (σ (*var-rule* α ! *i*)) $\sqcup$ *As*!*i* = *Some* (*Cs*!*i*))
using *join-rule-fun* **by** (*metis join-sym*)
from *5-2*(*4*) **have** *wellA*:$\forall\, i$ < *length As*. *As*!*i* $\in$ *wf-pterm R*
by *auto*
from *5-2*(*4*) **have** *l*: *length As* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
then have *l2*:*length As* = *length* (*zip As* (*map* σ (*var-rule* α)))
by *simp*
from *l* $*$ **have** $\forall\, i$ < *length As*. σ (*var-rule* α ! *i*) $\in$ *wf-pterm R*
using *5-2*(*3*) **by** (*metis match-well-def vars-to-pterm*)
with *5-2*(*1*) $*$ *wellA l2 l* **have** *IH*:$\forall\, i$ < *length As*. *co-initial* ((*map* σ (*var-rule* α))!*i*) (*Cs*!*i*)
by (*smt* (*verit*, *del-insts*) *length-map nth-map nth-mem nth-zip zip-symm*)
moreover have *f*:*Pfun f Bs* = (*to-pterm* (*lhs* α)) $\cdot$ $\langle$(*map* σ (*var-rule* α))$\rangle_\alpha$
using $*$ **by** (*smt* (*verit*, *ccfv-threshold*) *apply-lhs-subst-var-rule map-eq-conv match-lhs-subst*)
show *?case* **using** *IH l* **unfolding** *f* $*$(*2*) *source.simps*
by (*metis* $*$(*3*) *fun-mk-subst length-map nth-map-conv source.simps*(*1*) *source-apply-subst source-to-pterm to-pterm-wf-pterm*)
qed *auto*

lemma *join-pterm-subst-Some*:
fixes *A B*::('*f*, '*v*) *pterm*
assumes (*A* $\cdot$ σ) $\sqcup$ (*A* $\cdot$ τ) = *Some B*
shows $\exists\, \varrho$. ($\forall\, x \in$ *vars-term A*. σ *x* $\sqcup$ τ *x* = *Some* (ϱ *x*)) $\wedge$ *B* = *A* $\cdot$ ϱ $\wedge$ *match B A* = *Some* ϱ
proof−
let *?join-var*=λx. *the* (σ *x* $\sqcup$ τ *x*)
let *?*ϱ=*mk-subst Var* (*zip* (*vars-distinct A*) (*map ?join-var* (*vars-distinct A*)))
from *assms* **have** *B* = *A* $\cdot$ *?*ϱ $\wedge$ ($\forall\, x \in$ *vars-term A*. σ *x* $\sqcup$ τ *x* = *Some* (*?*ϱ *x*)) $\wedge$ *match B A* = *Some ?*ϱ **proof**(*induct A arbitrary*: *B*)
case (*Var x*)
then show *?case*
by (*smt* (*verit*) *comp-apply eval-term.simps*(*1*) *in-set-conv-nth in-set-simps*(*2*) *length-map map-nth-conv match-trivial mem-Collect-eq mk-subst-not-mem mk-subst-same option.sel remdups-id-iff-distinct set-vars-term-list single-var singleton-rev-conv subsetI subst-domain-def vars-term-list.simps*(*1*))
next
case (*Pfun f As*)
let *?*ϱ=*mk-subst Var* (*zip* (*vars-distinct* (*Pfun f As*)) (*map ?join-var* (*vars-distinct* (*Pfun f As*))))
have *rho-domain*:*subst-domain ?*ϱ $\subseteq$ *vars-term* (*Pfun f As*)
by (*smt* (*verit*, *del-insts*) *comp-apply mem-Collect-eq mk-subst-not-mem*

set-remdups set-rev set-vars-term-list subsetI subst-domain-def)
from *Pfun*(*2*) **have** *those* (*map2* (⊔) (*map* (λ*a*. *a* · σ) *As*) (*map* (λ*a*. *a* · τ) *As*)) ≠ *None*
unfolding *eval-term.simps join.simps* **using** *option.simps*(*4*) **by** *fastforce*
then obtain *Bs* **where** *Bs*:*those* (*map2* (⊔) (*map* (λ*a*. *a* · σ) *As*) (*map* (λ*a*. *a* · τ) *As*)) = *Some Bs length As* = *length Bs*
using *length-those* **by** *fastforce*
{**fix** *i* **assume** *i* < *length As*
with *Bs* **have** *Bs-i*:((*As*!*i*) · σ) ⊔ ((*As*!*i*) · τ) = *Some* (*Bs*!*i*)
using *those-some2* **by** *fastforce*
}**note** *Bs-i*=*this*
{**fix** *i* **assume** *i*:*i* < *length As*
let *?ϱi*=*mk-subst Var* (*zip* (*vars-distinct* (*As*!*i*)) (*map ?join-var* (*vars-distinct* (*As*!*i*))))
have (*As*!*i*) · *?ϱ* = (*As*!*i*) · *?ϱi*
by (*smt* (*verit*, *ccfv-SIG*) *comp-apply i map-of-zip-map mk-subst-def nth-mem set-remdups set-rev set-vars-term-list term.set-intros*(*4*) *term-subst-eq-conv*)
with *Pfun*(*1*)[*of As*!*i*] *i Bs-i* **have** (*As*!*i*) · *?ϱ* = *Bs*!*i*
by *fastforce*
}**note** *As-Bs*=*this*
with *Bs*(*2*) **have** *map-ϱ*:*map* (λ*a*. *a* · *?ϱ*) *As* = *Bs*
by (*simp add*: *map-nth-eq-conv*)
{**fix** *x* **assume** *x*:*x* ∈ *vars-term* (*Pfun f As*)
then obtain *i* **where** *i* < *length As x* ∈ *vars-term* (*As*!*i*)
by (*metis term.sel*(*4*) *var-imp-var-of-arg*)
with *Pfun*(*1*)[*of As*!*i*] *Bs-i As-Bs* **have** σ *x* ⊔ τ *x* = *Some* (*?ϱ x*)
using *term-subst-eq-rev* **by** *fastforce*
}
moreover then have *B* = *Pfun f As* · *?ϱ*
using *Pfun*(*2*) **unfolding** *eval-term.simps join.simps Bs* **using** *map-ϱ* **by** *auto*
moreover then have *match B* (*Pfun f As*) = *Some ?ϱ*
using *match-trivial rho-domain* **by** *blast*
ultimately show *?case* **by** *simp*
next
case (*Prule* α *As*)
let *?ϱ*=*mk-subst Var* (*zip* (*vars-distinct* (*Prule* α *As*)) (*map ?join-var* (*vars-distinct* (*Prule* α *As*))))
have *rho-domain*:*subst-domain ?ϱ* ⊆ *vars-term* (*Prule* α *As*)
by (*smt* (*verit*, *del-insts*) *comp-apply mem-Collect-eq mk-subst-not-mem set-remdups set-rev set-vars-term-list subsetI subst-domain-def*)
from *Prule*(*2*) **have** *those* (*map2* (⊔) (*map* (λ*a*. *a* · σ) *As*) (*map* (λ*a*. *a* · τ) *As*)) ≠ *None*
unfolding *eval-term.simps join.simps* **using** *option.simps*(*4*) **by** *fastforce*
then obtain *Bs* **where** *Bs*:*those* (*map2* (⊔) (*map* (λ*a*. *a* · σ) *As*) (*map* (λ*a*. *a* · τ) *As*)) = *Some Bs length As* = *length Bs*
using *length-those* **by** *fastforce*
{**fix** *i* **assume** *i* < *length As*
with *Bs* **have** *Bs-i*:((*As*!*i*) · σ) ⊔ ((*As*!*i*) · τ) = *Some* (*Bs*!*i*)

using *those-some2* **by** *fastforce*
}note *Bs-i=this*
{fix *i* **assume** *i*:*i* < *length As*
let *?ϱi=mk-subst Var* (*zip* (*vars-distinct* (*As*!*i*)) (*map ?join-var* (*vars-distinct* (*As*!*i*))))
have (*As*!*i*) · *?ϱ* = (*As*!*i*) · *?ϱi*
by (*smt* (*verit*, *ccfv-SIG*) *comp-apply i map-of-zip-map mk-subst-def nth-mem set-remdups set-rev set-vars-term-list term.set-intros*(*4*) *term-subst-eq-conv*)
with *Prule*(*1*)[*of As*!*i*] *i Bs-i* **have** (*As*!*i*) · *?ϱ* = *Bs*!*i*
by *fastforce*
}note *As-Bs=this*
with *Bs*(*2*) **have** *map-ϱ*:*map* (*λa*. *a* · *?ϱ*) *As* = *Bs*
by (*simp add*: *map-nth-eq-conv*)
{fix *x* **assume** *x*:*x* ∈ *vars-term* (*Prule α As*)
then obtain *i* **where** *i* < *length As x* ∈ *vars-term* (*As*!*i*)
by (*metis term.sel*(*4*) *var-imp-var-of-arg*)
with *Prule*(*1*)[*of As*!*i*] *Bs-i As-Bs* **have** *σ x* ⊔ *τ x* = *Some* (*?ϱ x*)
using *term-subst-eq-rev* **by** *fastforce*
}
moreover then have *B* = *Prule α As* · *?ϱ*
using *Prule*(*2*) **unfolding** *eval-term.simps join.simps Bs* **using** *map-ϱ* **by** *auto*
moreover then have *match B* (*Prule α As*) = *Some ?ϱ*
using *match-trivial rho-domain* **by** *blast*
ultimately show *?case* **by** *simp*
qed
then show *?thesis* **by** *blast*
qed

lemma *join-pterm-subst-None*:
fixes *A*::(*'f*, *'v*) *pterm*
assumes (*A* · *σ*) ⊔ (*A* · *τ*) = *None*
shows ∃ *x* ∈ *vars-term A*. *σ x* ⊔ *τ x* = *None*
using *assms* **proof**(*induct A rule*:*pterm-induct*)
case (*Pfun f As*)
from *Pfun*(*2*) **obtain** *i* **where** *i*:*i* < *length As* (*map2* (⊔) (*map* (*λs*. *s* · *σ*) *As*) (*map* (*λs*. *s* · *τ*) *As*))!*i* = *None*
unfolding *eval-term.simps join.simps length-map* **using** *those-not-none-xs*
by (*smt* (*verit*) *length-map list-all-length map2-map-map option.case-eq-if option.distinct*(*1*))
then have (*As*!*i* · *σ*) ⊔ (*As*!*i* · *τ*) = *None* **by** *simp*
with *Pfun*(*1*) *i*(*1*) **obtain** *x* **where** *x* ∈ *vars-term* (*As*!*i*) **and** *σ x* ⊔ *τ x* = *None*
using *nth-mem* **by** *blast*
then show *?case* **using** *i*(*1*) **by** *auto*
next
case (*Prule α As*)
from *Prule*(*2*) **obtain** *i* **where** *i*:*i* < *length As* (*map2* (⊔) (*map* (*λs*. *s* · *σ*) *As*) (*map* (*λs*. *s* · *τ*) *As*))!*i* = *None*
unfolding *eval-term.simps join.simps length-map* **using** *those-not-none-xs*

by (*smt* (*verit*) *length-map list-all-length map2-map-map option.case-eq-if option.distinct*(*1*))
then have $(As!i \cdot \sigma) \sqcup (As!i \cdot \tau) = None$ **by** *simp*
with *Prule*(*1*) *i*(*1*) **obtain** x **where** $x \in$ *vars-term* $(As!i)$ **and** $\sigma\ x \sqcup \tau\ x = None$
using *nth-mem* **by** *blast*
then show *?case* **using** *i*(*1*) **by** *auto*
qed *simp*

fun *mk-subst-from-list* :: $('v \Rightarrow ('f, 'v)\ term)\ list \Rightarrow ('v \Rightarrow ('f, 'v)\ term)$ **where**
mk-subst-from-list [] = *Var*
| *mk-subst-from-list* $(\sigma \# \sigma s) = (\lambda x.\ case\ \sigma\ x\ of$
$Var\ x \Rightarrow$ *mk-subst-from-list* $\sigma s\ x$
$|\ t \Rightarrow t)$

lemma *join-is-Fun*:
assumes *join*:$A \sqcup B = Some\ (Pfun\ f\ Cs)$
shows $\exists As.\ A = Pfun\ f\ As \wedge length\ As = length\ Cs$
proof−
{**assume** $\exists x.\ A = Var\ x$
then obtain x **where** $x{:}A = Var\ x$ **by** *blast*
from *join* **consider** $B = Var\ x \mid \exists \alpha\ Bs.\ B = Prule\ \alpha\ Bs$
unfolding x **by** (*metis is-Prule.elims*(*1*) *join.simps*(*1*) *join.simps*(*9*) *option.distinct*(*1*))
then have *False*
using *join* **unfolding** x **by**(*cases*) (*simp*, *metis* (*mono-tags*, *lifting*) *Inl-Inr-False join.simps*(*6*) *option.case-eq-if option.distinct*(*1*) *option.sel term.inject*(*2*))
} **moreover** {**assume** $\exists \alpha\ As.\ A = Prule\ \alpha\ As$
then obtain $\alpha\ As$ **where** $A{:}A = Prule\ \alpha\ As$ **by** *blast*
from *join* **consider** $\exists x.\ B = Var\ x \mid \exists g\ Bs.\ B = Pfun\ g\ Bs$
unfolding A **by** (*smt* (*verit*, *del-insts*) *is-Prule.elims*(*1*) *join.simps*(*3*) *option.case-eq-if option.distinct*(*1*) *option.sel term.inject*(*2*))
then have *False*
using *join* **unfolding** A **by**(*cases*) (*metis* (*mono-tags*, *lifting*) *Inl-Inr-False join.simps*(*4*,*5*) *option.case-eq-if option.distinct*(*1*) *option.sel term.inject*(*2*))+
}
ultimately obtain $g\ As$ **where** $A{:}A = Pfun\ g\ As$
by (*meson is-Prule.cases*)
from *join* **have** $f = g$ **and** $length\ As = length\ Cs$ **unfolding** A
by (*smt* (*verit*, *ccfv-threshold*) *Inl-Inr-False Residual-Join-Deletion.join-sym join.simps*(*5*) *join.simps*(*8*) *join-fun-fun not-arg-cong-Inr option.case-eq-if option.inject option.simps*(*3*) *pterm-cases term.inject*(*2*))+
with A **show** *?thesis* **by** *force*
qed

lemma *join-obtain-subst*:
assumes *join*:$A \sqcup B = Some\ (to\text{-}pterm\ t \cdot \sigma)$ **and** *linear-term* t
shows $(to\text{-}pterm\ t) \cdot mk\text{-}subst\ Var\ (match\text{-}substs\ (to\text{-}pterm\ t)\ A) = A$
proof−

from *assms*(*2*) **have** *lin*:*linear-term* (*to-pterm t*)
using *to-pterm-linear* **by** *blast*
have $\forall$ *p*$\in$*poss* (*to-pterm t*). $\forall$ *f ts*. (*to-pterm t*) |- *p* = *Fun f ts* $\longrightarrow$ ($\exists$ *As*. *length ts* = *length As* $\wedge$ *A* |- *p* = *Fun f As*)
using *assms* **proof**(*induct t arbitrary*: *A B*)
case (*Fun f ts*)
from *Fun*(*2*) **obtain** *As* **where** *A*:*A* = *Pfun f As* **and** *l-As*:*length ts* = *length As*
using *join-is-Fun* **by** *force*
from *Fun*(*2*) **obtain** *Bs* **where** *B*:*B* = *Pfun f Bs* **and** *l-Bs*:*length ts* = *length Bs*
using *join-is-Fun join-sym* **by** (*smt* (*verit*) *eval-term.simps*(*2*) *length-map to-pterm.simps*(*2*))
{**fix** *p g ts*′ **assume** *p*:*p* $\in$ *poss* (*to-pterm* (*Fun f ts*)) (*to-pterm* (*Fun f ts*)) |-*p* = *Fun g ts*′
have $\exists$ *As*′. *length ts*′ = *length As*′ $\wedge$ *A*|-*p* = *Fun g As*′ **proof**(*cases p*)
case *Nil*
from *p*(*2*) **show** *?thesis* **unfolding** *A Nil* **using** *l-As* **by** *force*
next
case (*Cons i p*′)
from *p*(*1*) **have** *i*:*i* < *length ts* **unfolding** *Cons* **by** *simp*
with *p*(*1*) **have** *p*′:*p*′ $\in$ *poss* (*ts*!*i*) **unfolding** *Cons*
by (*metis poss-Cons-poss poss-list-sound poss-list-to-pterm term.sel*(*4*))
from *Fun*(*2*) **have** *As*!*i* $\sqcup$ *Bs*!*i* = *Some* (*to-pterm* (*ts*!*i*) $\cdot$ σ)
unfolding *A B to-pterm.simps eval-term.simps* **using** *i l-As l-Bs*
by (*smt* (*verit*, *ccfv-threshold*) *args-poss join-fun-fun local.Cons nth-map p*(*1*) *term.sel*(*4*) *to-pterm.simps*(*2*))
moreover from *Fun*(*3*) *i* **have** *linear-term* (*ts*!*i*) **by** *simp*
ultimately obtain *As*′ **where** *length ts*′ = *length As*′ **and** (*As*!*i*)|-*p*′ = *Fun g As*′
using *Fun*(*1*) *i p*′ **by** (*smt* (*verit*) *local.Cons nth-map nth-mem p*(*2*) *p-in-poss-to-pterm subt-at.simps*(*2*) *to-pterm.simps*(*2*))
with *i l-As* **show** *?thesis* **unfolding** *A Cons* **by** *simp*
qed
}
then show *?case* **by** *simp*
qed *simp*
then show *?thesis* **using** *fun-poss-eq-imp-matches*[*OF lin*] **by** *presburger*
qed

lemma *join-pterm-linear-subst*:
assumes *join*:*A* $\sqcup$ *B* = *Some* (*to-pterm t* $\cdot$ σ) **and** *lin*:*linear-term t*
shows $\exists$ σ_A σ_B. *A* = (*to-pterm t* $\cdot$ σ_A) $\wedge$ *B* = (*to-pterm t* $\cdot$ σ_B) $\wedge$ ($\forall$ *x* $\in$ *vars-term t*. σ_A *x* $\sqcup$ σ_B *x* = *Some* (σ *x*))
proof−
let *?*σ_A=*mk-subst Var* (*match-substs* (*to-pterm t*) *A*)
let *?*σ_B=*mk-subst Var* (*match-substs* (*to-pterm t*) *B*)
from *join-obtain-subst*[*OF join lin*] **have** *A*:*A* = (*to-pterm t*) $\cdot$ *?*σ_A **by** *simp*
from *join lin* **have** *B*:*B* = (*to-pterm t*) $\cdot$ *?*σ_B **using** *join-sym join-obtain-subst*

by *metis*
  from *join-pterm-subst-Some join A B* **obtain** ϱ
    where $(\forall x \in vars\text{-}term\ t.\ ?\sigma_A\ x \sqcup ?\sigma_B\ x = Some\ (\varrho\ x))$ **and** $to\text{-}pterm\ t \cdot \sigma = to\text{-}pterm\ t \cdot \varrho$
    by (*metis set-vars-term-list vars-to-pterm*)
  then show *?thesis*
    by (*smt* (*verit*, *best*) *A B set-vars-term-list term-subst-eq-rev vars-to-pterm*)
qed

context *no-var-lhs*
begin
lemma *join-rule-lhs*:
  assumes *wf*:$Prule\ \alpha\ As \in wf\text{-}pterm\ R$ **and** *args*:$\forall i < length\ As.\ As!i \sqcup Bs!i \neq None$ **and** *l*:$length\ Bs = length\ As$
  shows $Prule\ \alpha\ As \sqcup (to\text{-}pterm\ (lhs\ \alpha) \cdot \langle Bs\rangle_\alpha) \neq None$
proof−
  from *wf no-var-lhs* **obtain** *f ts* **where** *lhs*:$lhs\ \alpha = Fun\ f\ ts$
   by (*metis Inl-inject Term.term.simps*(*2*) *Term.term.simps*(*4*) *case-prodD is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *term.collapse*(*2*) *wf-pterm.simps*)
  from *args l* **have** $those\ (map2\ (\sqcup)\ As\ Bs) \neq None$
    by (*simp add*: *list-all-length those-not-none-xs*)
  with *wf l* **have** $those\ (map2\ (\sqcup)\ As\ (map\ (\langle Bs\rangle_\alpha)\ (vars\text{-}distinct\ (Fun\ f\ ts)))) \neq None$
   using *apply-lhs-subst-var-rule* **by** (*metis Inl-inject is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *lhs term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.simps*)
  with *lhs-subst-trivial*[*of* α *Bs*] **show** *?thesis*
    unfolding *lhs to-pterm.simps eval-term.simps join.simps* **by** *force*
qed
end

3.2.1 N-Fold Join

We define a function to recursively join a list of n proof terms. Since each individual join produces a $(('f, 'v)\ prule + 'f, 'v)\ Term.term\ option$ we first introduce the following helper function.

fun *join-opt* :: $('f, 'v)\ pterm \Rightarrow ('f, 'v)\ pterm\ option \Rightarrow ('f, 'v)\ pterm\ option$
  where
    $join\text{-}opt\ A\ (Some\ B) = A \sqcup B$
  | $join\text{-}opt\ \text{-}\ \text{-} = None$

fun *join-list* :: $('f, 'v)\ pterm\ list \Rightarrow ('f,'v)\ pterm\ option\ (\bigsqcup)$
  where
    $join\text{-}list\ [] = None$
  | $join\text{-}list\ (A\ \#\ []) = Some\ A$
  | $join\text{-}list\ (A\ \#\ As) = join\text{-}opt\ A\ (join\text{-}list\ As)$

context *left-lin-no-var-lhs*
begin

lemma *join-var-rule*:
  assumes *to-rule* $\alpha \in R$
  shows *Var x* $\sqcup$ *Prule* α *As* = *None*
proof –
  from *assms* **obtain** *f ts* **where** *lhs* α = *Fun f ts*
    using *no-var-lhs* **by** *fastforce*
  then show *?thesis*
   by (*metis* (*no-types*, *lifting*) *Residual-Join-Deletion.join-sym eval-term.simps*(*2*)
join.simps(*4*) *match-lhs-subst option.case-eq-if option.exhaust term.distinct*(*1*) *to-pterm.simps*(*2*))
qed

lemma *var-join*:
  assumes *Var x* $\sqcup$ *B* = *Some C* **and** *B* $\in$ *wf-pterm R*
  shows *B* = *Var x* $\wedge$ *C* = *Var x*
  using *assms* **proof**(*cases B*)
  case (*Var y*)
  with *assms*(*1*) **show** *?thesis*
    by (*metis join.simps*(*1*) *option.sel option.simps*(*3*))
next
  case (*Prule* α *As*)
  with *assms* **show** *?thesis*
      by (*metis Residual-Join-Deletion.join-sym Term.term.simps*(*4*) *case-prodD*
co-initial-Var is-VarI join.simps(*9*)
      no-var-lhs.no-var-lhs no-var-lhs-axioms option.distinct(*1*) *source-join sum.inject*(*1*)
term.inject(*2*) *wf-pterm.simps*)
qed *simp*

lemma *fun-join*:
  assumes *Pfun f As* $\sqcup$ *B* = *Some C*
  shows ($\exists$ *g Bs*. *B* = *Pfun g Bs*) $\vee$ ($\exists$ α *Bs*. *B* = *Prule* α *Bs*)
  using *assms* **by**(*cases B*) (*simp-all*)

lemma *rule-join*:
  assumes *Prule* α *As* $\sqcup$ *B* = *Some C* **and** *Prule* α *As* $\in$ *wf-pterm R*
  shows ($\exists$ *g Bs*. *B* = *Pfun g Bs*) $\vee$ ($\exists$ β *Bs*. *B* = *Prule* β *Bs*)
  using *assms* **proof**(*cases B*)
  case (*Var x*)
  from *assms* **have** *False* **unfolding** *Var*
    by (*metis Residual-Join-Deletion.join-sym term.distinct*(*1*) *var-join*)
  then show *?thesis* **by** *simp*
qed *simp-all*

Associativity of join is currently not used in any proofs. But it is still a valuable result, hence included here.

lemma *join-opt-assoc*:
  assumes *A* $\in$ *wf-pterm R B* $\in$ *wf-pterm R C* $\in$ *wf-pterm R*
  shows *join-opt A* (*B* $\sqcup$ *C*) = *join-opt C* (*A* $\sqcup$ *B*)
  using *assms* **proof**(*induct A arbitrary*:*B C rule*:*subterm-induct*)
  case (*subterm A*)

from *subterm(2)* **show** *?case* **proof**(*cases A rule*:*wf-pterm.cases*[*case-names VarA FunA RuleA*])
case (*VarA x*)
with *subterm(3)* **show** *?thesis* **proof**(*cases B rule*:*wf-pterm.cases*[*case-names VarB FunB RuleB*])
case (*VarB y*)
show *?thesis* **proof**(*cases* $x = y$)
case *True*
then have *AB*:$A \sqcup B$ = *Some* (*Var y*) **unfolding** *VarA VarB* **by** *simp*
with *subterm(4) VarA VarB* **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC z*)
with *AB VarA VarB* ‹$x = y$› **show** *?thesis* **by**(*cases* $y = z$) *simp-all*
next
case (*RuleC* α *As*)
then have (*Var y*) $\sqcup$ *C* = *None*
using *join-var-rule* **by** *presburger*
then show *?thesis* **unfolding** *AB* **unfolding** *VarA VarB* **by** (*simp add*:*join-sym*)
qed *simp*
next
case *False*
then have *AB*:$A \sqcup B$ = *None* **unfolding** *VarA VarB* **by** *simp*
with *subterm(4) VarA VarB* **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC z*)
with *AB VarA VarB* ‹$x \neq y$› **show** *?thesis* **by**(*cases* $y = z$) *simp-all*
next
case (*RuleC* α *As*)
then have (*Var y*) $\sqcup$ *C* = *None*
using *join-var-rule* **by** *presburger*
then show *?thesis* **unfolding** *AB* **unfolding** *VarA VarB* **by** (*simp add*:*join-sym*)
qed *simp*
qed
next
case (*FunB Bs f*)
then have *AB*:$A \sqcup B$ = *None*
unfolding *VarA* **by** *simp*
with *subterm(4) VarA* **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC z*)
with *AB VarA FunB* **show** *?thesis* **by**(*cases* $x = z$) *simp-all*
next
case (*FunC Cs g*)
from *AB VarA* **show** *?thesis* **proof**(*cases* $B \sqcup C$)
case (*Some BC*)
then obtain *BCs* **where** *BC* = *Pfun f BCs*
by (*metis FunB(1) FunC(1) join-fun-fun*)

then show *?thesis* **unfolding** *AB* **unfolding** *VarA Some* **by** *simp*
qed *simp*
next
case $(RuleC\ \alpha\ Cs)$
from *AB VarA* **show** *?thesis* **proof**(*cases* $B \sqcup C$)
case (*Some BC*)
then obtain *BCs* **where** $BC = Prule\ \alpha\ BCs$
by (*metis FunB*(*1*) *Residual-Join-Deletion.join-sym RuleC*(*1*) *join-rule-fun subterm.prems*(*3*))
then have $Var\ x \sqcup BC = None$
using *RuleC*(*2*) *join-var-rule* **by** *presburger*
then show *?thesis* **unfolding** *AB* **unfolding** *VarA Some* **by** *simp*
qed *simp*
qed
next
case $(RuleB\ \alpha\ Bs)$
then have $AB{:}A \sqcup B = None$
using *VarA join-var-rule* **by** *presburger*
with *subterm*(*4*) *VarA* **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC z*)
with *RuleB* **have** $B \sqcup C = None$
using *join-var-rule join-sym* **by** *metis*
with *AB* **show** *?thesis* **by** *simp*
next
case (*FunC Cs f*)
from *AB VarA* **show** *?thesis* **proof**(*cases* $B \sqcup C$)
case (*Some BC*)
then obtain *BCs* **where** $BC = Prule\ \alpha\ BCs$
by (*metis FunC*(*1*) *RuleB*(*1*) *join-rule-fun subterm.prems*(*2*))
then have $Var\ x \sqcup BC = None$
using *RuleB*(*2*) *join-var-rule* **by** *presburger*
then show *?thesis* **unfolding** *AB* **unfolding** *VarA Some* **by** *simp*
qed *simp*
next
case $(RuleC\ \beta\ Cs)$
from *AB VarA* **show** *?thesis* **proof**(*cases* $B \sqcup C$)
case (*Some BC*)
then obtain *BCs* **where** $BC = Prule\ \alpha\ BCs$
using *RuleB*(*1*) *RuleC*(*1*) *join-rule-rule subterm.prems*(*2*) *subterm.prems*(*3*) **by** *blast*
then have $Var\ x \sqcup BC = None$
using *RuleB*(*2*) *join-var-rule* **by** *presburger*
then show *?thesis* **unfolding** *AB* **unfolding** *VarA Some* **by** *simp*
qed *simp*
qed
qed
next
case (*FunA As f*)

from *subterm*(*3*) **show** *?thesis* **proof**(*cases B rule*:*wf-pterm.cases*[*case-names VarB FunB RuleB*])
case (*VarB x*)
then show *?thesis*
by (*metis FunA*(*1*) *join.simps*(*1*) *join.simps*(*8*) *join.simps*(*9*) *join-opt.elims join-opt.simps*(*2*) *join-var-rule option.sel subterm.prems*(*3*) *wf-pterm.simps*)
next
case (*FunB Bs g*)
then show *?thesis* **proof**(*cases* $A \sqcup B$)
case *None*
with *subterm*(*4*) *FunB* **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*FunC Cs h*)
from *None* **show** *?thesis* **proof**(*cases* $B \sqcup C$)
case (*Some BC*)
then have *gh*:$g = h$ **and** *l-B-C*:*length Bs* = *length Cs*
unfolding *FunB FunC* **by** (*meson join-fun-fun*)+
from *Some* **obtain** *BCs* **where** *BC*:*BC* = *Pfun g BCs* **and** *l-BC-B*:*length BCs* = *length Bs*
and *args-BC*:($\forall$ *i*<*length Bs*. *Bs*!*i* $\sqcup$ *Cs*!*i* = *Some* (*BCs* ! *i*))
unfolding *FunC FunB* **using** *join-fun-fun* **by** *force*
{**assume** *fg*:$f = g$ **and** *l-A-B*:*length As* = *length Bs*
{**fix** *i* **assume** *i*:*i* < *length As*
with *subterm*(*1*) *FunA FunB*(*2*) *FunC*(*2*) *args-BC l-A-B l-B-C*
have *join-opt* (*As*!*i*) ((*Bs*!*i*) $\sqcup$ (*Cs*!*i*)) = *join-opt* (*Cs*!*i*) ((*As*!*i*) $\sqcup$ (*Bs*!*i*))
by (*metis nth-mem supt.arg*)
} **note** *IH*=*this*
from *fg l-A-B None* **have** *those* (*map2* ($\sqcup$) *As Bs*) = *None*
unfolding *FunB FunA* **by** (*smt* (*verit*) *join.simps*(*2*) *option.case-eq-if option.distinct*(*1*))
then obtain *i* **where** *i*:*i* < *length* (*map2* ($\sqcup$) *As Bs*) (*map2* ($\sqcup$) *As Bs*)!*i* = *None*
using *those-not-none-xs list-all-length* **by** *blast*
with *l-A-B* **have** *A-B-i*:(*As*!*i*) $\sqcup$ (*Bs*!*i*) = *None* **by** *simp*
with *IH i*(*1*) *l-B-C* **have** *As*!*i* $\sqcup$ *BCs*!*i* = *None* **using** *args-BC* **by** *fastforce*
with *i*(*1*) *l-BC-B l-B-C l-A-B* **have** *those* (*map2* ($\sqcup$) *As BCs*) = *None*
using *list-all-length those-some2* **by** *fastforce*
}
then show *?thesis*
using *l-B-C l-BC-B FunA FunB FunC BC gh None Some* **by** *auto*
qed *simp*
next
case (*RuleC* α *Cs*)
from *None* **show** *?thesis* **proof**(*cases* $B \sqcup C$)
case (*Some BC*)
then obtain *BCs* τ **where** τ:*match B* (*to-pterm* (*lhs* α)) = *Some* τ **and** *BC*:*BC* = *Prule* α *BCs*

and *l-BCs*:*length BCs* = *length Cs* **and** *args-BC*:$\forall$ *i* < *length Cs*. *Cs*!*i* $\sqcup$ τ (*var-rule* α ! *i*) = *Some* (*BCs* ! *i*)
by (*metis FunB*(*1*) *join-sym RuleC*(*1*) *join-rule-fun subterm.prems*(*3*))

with *None Some FunA* **show** *?thesis* **proof**(*cases match A* (*to-pterm* (*lhs* α)))
case (*Some* σ)
with τ *None* **obtain** *x* **where** *x*:*x* $\in$ *vars-term* (*lhs* α) σ *x* $\sqcup$ τ *x* = *None*
using *join-pterm-subst-None* **by** (*metis lhs-subst-trivial match-lhs-subst option.sel set-vars-term-list vars-to-pterm*)
then obtain *i* **where** *i*:*i* < *length* (*var-rule* α) *var-rule* α ! *i* = *x*
by (*metis RuleC*(*2*) *case-prodD in-set-idx left-lin left-linear-trs-def linear-term-var-vars-term-list set-vars-term-list*)
have *subt*:*A* $\rhd$ σ *x* **proof**−
obtain *g ts* **where** *lhs*:*lhs* α = *Fun g ts*
using *RuleC*(*2*) *no-var-lhs* **by** *fastforce*
from *Some i* **show** *?thesis*
unfolding *lhs* **by** (*metis* (*no-types*, *lifting*) *lhs match-matches set-vars-term-list subst-image-subterm to-pterm.simps*(*2*) *vars-to-pterm x*(*1*))
qed
have *wf-*τ*-x*:τ *x* $\in$ *wf-pterm R*
using *FunB* τ *i* **by** (*metis match-well-def subterm.prems*(*2*) *vars-to-pterm*)
have *IH*:*join-opt* (σ *x*) (τ *x* $\sqcup$ (*Cs* ! *i*)) = *join-opt* (*Cs*!*i*) (σ *x* $\sqcup$ τ *x*)
using *subterm*(*1*) *RuleC*(*3*,*4*) *i wf-*τ*-x* **by** (*metis Some match-well-def nth-mem subt subterm.prems*(*1*) *vars-to-pterm*)
have τ *x* $\sqcup$ (*Cs* ! *i*) = *Some* (*BCs* ! *i*)
using *args-BC i* **by** (*metis Residual-Join-Deletion.join-sym RuleC*(*3*))

with *IH x*(*2*) **have** (σ *x*) $\sqcup$ (*BCs* ! *i*) = *None* **by** *simp*
then have (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *BCs*) ! *i* = *None*
using *l-BCs i* **by** (*simp add*: *RuleC*(*3*))
then have *those* (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *BCs*) = *None*
using *l-BCs i* **by** (*metis* (*no-types*, *opaque-lifting*) *RuleC*(*3*) *length-map length-zip min.idem not-Some-eq nth-mem those-not-none-x*)
then have *A* $\sqcup$ *BC* = *None*
using *Some i* **unfolding** *BC FunA join.simps* **by** *simp*
then show *?thesis*
unfolding *None* ‹*B* $\sqcup$ *C* = *Some BC*› **by** *auto*
qed *simp*
qed *simp*
qed *simp*
next
case (*Some AB*)
then have *fg*:*f* = *g* **and** *l-A-B*:*length As* = *length Bs*
unfolding *FunA FunB* **by** (*meson join-fun-fun*)+
from *Some* **obtain** *ABs* **where** *AB*:*AB* = *Pfun f ABs* **and** *l-AB-A*:*length ABs* = *length As*

and *args-AB*:(∀ *i*<*length Bs*. *As*!*i* ⊔ *Bs*!*i* = *Some* (*ABs* ! *i*))
unfolding *FunA FunB* **using** *join-fun-fun* **by** *force*
from *subterm*(*4*) **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC x*)
show *?thesis* **unfolding** *Some AB* **unfolding** *FunA FunB VarC* **by** *simp*
next
case (*FunC Cs h*)
show *?thesis* **proof**(*cases B* ⊔ *C*)
case *None*
{**assume** *gh*:*g* = *h* **and** *l-B-C*:*length Bs* = *length Cs*
{**fix** *i* **assume** *i*:*i* < *length As*
with *subterm*(*1*) *FunA FunB*(*2*) *FunC*(*2*) *args-AB l-A-B l-B-C*
have *join-opt* (*As*!*i*) ((*Bs*!*i*) ⊔ (*Cs*!*i*)) = *join-opt* (*Cs*!*i*) ((*As*!*i*) ⊔ (*Bs*!*i*))
by (*metis nth-mem supt.arg*)
} **note** *IH*=*this*
from *gh l-B-C None* **have** *those* (*map2* (⊔) *Bs Cs*) = *None*
unfolding *FunB FunC* **by** (*smt* (*verit*) *join.simps*(*2*) *option.case-eq-if option.distinct*(*1*))
then obtain *i* **where** *i*:*i* < *length* (*map2* (⊔) *Bs Cs*) (*map2* (⊔) *Bs Cs*)!*i* = *None*
using *those-not-none-xs list-all-length* **by** *blast*
with *l-B-C* **have** *B-C-i*:(*Bs*!*i*) ⊔ (*Cs*!*i*) = *None* **by** *simp*
with *IH i*(*1*) *l-A-B* **have** *Cs*!*i* ⊔ *ABs*!*i* = *None* **using** *args-AB* **by** *fastforce*
with *i*(*1*) *l-AB-A l-B-C l-A-B* **have** *those* (*map2* (⊔) *Cs ABs*) = *None*
using *list-all-length those-some2* **by** *fastforce*
}
then show *?thesis*
using *l-A-B l-AB-A FunA FunB FunC AB fg None Some* **by** *auto*
next
case (*Some BC*)
then have *gh*:*g* = *h* **and** *l-B-C*:*length Bs* = *length Cs*
unfolding *FunB FunC* **by**(*meson join-fun-fun*)+
from *Some* **obtain** *BCs* **where** *BC*:*BC* = *Pfun g BCs* **and** *l-BC-B*:*length BCs* = *length Bs*
and *args-BC*:(∀ *i*<*length Cs*. *Bs*!*i* ⊔ *Cs*!*i* = *Some* (*BCs* ! *i*))
unfolding *FunB FunC* **using** *join-fun-fun* **by** *force*
{**fix** *i* **assume** *i*:*i* < *length As*
with *subterm*(*1*) *FunA FunB*(*2*) *FunC*(*2*) *args-AB l-A-B l-B-C*
have *join-opt* (*As*!*i*) ((*Bs*!*i*) ⊔ (*Cs*!*i*)) = (*Cs*!*i*) ⊔ (*ABs*!*i*)
by (*metis join-opt.simps*(*1*) *nth-mem supt.arg*)
with *args-BC i l-A-B l-B-C* **have** (*As*!*i*) ⊔ (*BCs*!*i*) = (*Cs*!*i*) ⊔ (*ABs*!*i*) **by** *simp*
} **note** *IH*=*this*
then have *those* (*map2* (⊔) *As BCs*) = *those* (*map2* (⊔) *Cs ABs*)
by (*smt* (*verit*, *del-insts*) *l-AB-A l-A-B l-BC-B l-B-C length-zip map-equality-iff min-less-iff-conj nth-zip old.prod.case*)

then show *?thesis* **unfolding** *Some BC* ‹$A \sqcup B = \mathit{Some}\ AB$› *AB* **unfolding** *gh FunA FunC fg join-opt.simps* **using** *l-BC-B l-AB-A l-A-B l-B-C* **by** *simp*
qed
next
case (*RuleC* α *Cs*)
from *RuleC*(*2*) **have** *lin*:*linear-term* (*lhs* α)
using *left-lin left-linear-trs-def* **by** *fastforce*
from *RuleC*(*2*) **obtain** f' *ts* **where** *lhs*:*lhs* α = *Fun* f' *ts*
using *no-var-lhs* **by** *fastforce*
consider *match A* (*to-pterm* (*lhs* α)) = *None* | *match B* (*to-pterm* (*lhs* α)) = *None*
| (*matches*) *match A* (*to-pterm* (*lhs* α)) $\neq$ *None* $\wedge$ *match B* (*to-pterm* (*lhs* α)) $\neq$ *None* **by** *linarith*
then show *?thesis* **proof**(*cases*)
case *1*
then have *match*:*match AB* (*to-pterm* (*lhs* α)) = *None*
using *lin* **by** (*smt* (*verit*, *ccfv-threshold*) *Some join-pterm-linear-subst match-complete*$'$ *match-matches not-Some-eq*)
then have $C \sqcup AB = \mathit{None}$
unfolding *RuleC AB join.simps* **by** *simp*
moreover have *join-opt A* ($B \sqcup C$) = *None* **proof**−
consider ($\exists$ *BCs*. $B \sqcup C$ = *Some* (*Prule* α *BCs*)) | $B \sqcup C$ = *None*
unfolding *FunB RuleC join.simps* **by** (*metis* (*no-types*, *lifting*) *option.case-eq-if*)
then show *?thesis* **using** *1 FunA*(*1*) **by**(*cases*) (*force*, *simp*)
qed
ultimately show *?thesis* **using** *Some* **by** *simp*
next
case *2*
then have *match*:*match AB* (*to-pterm* (*lhs* α)) = *None*
using *lin* **by** (*smt* (*verit*, *ccfv-threshold*) *Some join-pterm-linear-subst match-complete*$'$ *match-matches not-Some-eq*)
then have $C \sqcup AB = \mathit{None}$
unfolding *RuleC AB join.simps* **by** *simp*
moreover from *2* **have** $B \sqcup C = \mathit{None}$
unfolding *FunB RuleC join.simps* **by** *simp*
ultimately show *?thesis* **using** *Some* **by** *simp*
next
case *matches*
from *matches* **obtain** σ **where** *sigma*:*match A* (*to-pterm* (*lhs* α)) = *Some* σ **by** *force*
from *matches* **obtain** τ **where** *tau*:*match B* (*to-pterm* (*lhs* α)) = *Some* τ **by** *force*
from *sigma tau* **obtain** ϱ **where** *rho*:($\forall x \in$ *vars-term* (*to-pterm* (*lhs* α)). $\sigma\ x \sqcup \tau\ x$ = *Some* ($\varrho\ x$))
and *AB-rho*:*AB* = (*to-pterm* (*lhs* α)) $\cdot\ \varrho$ **and** *match-rho*:*match AB* (*to-pterm* (*lhs* α)) = *Some* ϱ
using *join-pterm-subst-Some match-matches Some* **by** *blast*

{fix *i* **assume** *i*:$i <$ *length Cs*
with *sigma RuleC*(*3*) **have** (*map* σ (*var-rule* α))!$i \lhd A$
using *lhs* **by** (*smt* (*verit*, *ccfv-threshold*) *lin linear-term-var-vars-term-list match-matches nth-map nth-mem set-vars-term-list subst-image-subterm to-pterm.simps*(*2*) *vars-to-pterm*)
moreover have (*map* σ (*var-rule* α))!$i \in$ *wf-pterm R*
using *i match-well-def*[*OF subterm*(*2*) *sigma*] *RuleC*(*3*) **by** (*simp add*: *vars-to-pterm*)
moreover have (*map* τ (*var-rule* α))!$i \in$ *wf-pterm R*
using *i match-well-def*[*OF subterm*(*3*) *tau*] *RuleC*(*3*) **by** (*simp add*: *vars-to-pterm*)
ultimately have *join-opt* (*map* σ (*var-rule* α) ! *i*) (*map* τ (*var-rule* α) ! $i \sqcup$ *Cs*!*i*) = *Cs* ! $i \sqcup$ *map* ϱ (*var-rule* α) ! *i*
using *subterm*(*1*) *RuleC*(*3*,*4*) *i* **by** (*smt* (*verit*, *best*) *join-opt.simps*(*1*) *lin linear-term-var-vars-term-list nth-map nth-mem rho set-vars-term-list vars-to-pterm*)

}note *IH*=*this*
show *?thesis* **proof**(*cases those* (*map2* ($\sqcup$) (*map* τ (*var-rule* α)) *Cs*))
case *None*
then obtain *i* **where** *i*:$i <$ *length Cs* (*map* τ (*var-rule* α))!$i \sqcup$ *Cs*!*i* = *None*
using *those-not-none-xs* **by** (*smt* (*verit*) *length-map length-zip list-all-length map-nth-eq-conv min-less-iff-conj nth-zip old.prod.case*)
with *IH* **have** *Cs*!$i \sqcup$ *map* ϱ (*var-rule* α) ! *i* = *None* **by** *force*
with *i RuleC*(*3*) **have** $i <$ *length* (*map2* ($\sqcup$) *Cs* (*map* ϱ (*var-rule* α))) (*map2* ($\sqcup$) *Cs* (*map* ϱ (*var-rule* α))) ! *i* = *None* **by** *simp-all*
then have *those* (*map2* ($\sqcup$) *Cs* (*map* ϱ (*var-rule* α))) = *None*
by (*metis nth-mem option.exhaust those-not-none-x*)
with *None* **show** *?thesis* **unfolding** *Some* **unfolding** *FunA FunB RuleC join.simps tau*[*unfolded FunB*] **using** *AB match-rho* **by** *auto*
next
case (*Some BCs*)
then have *BC*:$B \sqcup C$ = *Some* (*Prule* α *BCs*)
unfolding *FunB RuleC join.simps tau*[*unfolded FunB*] **by** *simp*
from *Some* **have** *l-BCs*:*length BCs* = *length Cs*
using *RuleC*(*3*) *length-those* **by** *fastforce*
{fix *i* **assume** $i <$ *length Cs*
with *Some IH* **have** (*map* σ (*var-rule* α)) ! $i \sqcup$ *BCs* ! *i* = *Cs* ! $i \sqcup$ (*map* ϱ (*var-rule* α)) ! *i*
using *RuleC*(*3*) *those-some2* **by** *fastforce*
}
then have *map2* ($\sqcup$) (*map* σ (*var-rule* α)) *BCs* = *map2* ($\sqcup$) *Cs* (*map* ϱ (*var-rule* α))
using *l-BCs* **by** (*simp add*: *RuleC*(*3*) *map-eq-conv'*)
then show *?thesis* **unfolding** *BC* ‹$A \sqcup B$ = *Some AB*› **unfolding** *FunA FunB RuleC join-opt.simps join.simps sigma*[*unfolded FunA*] **using** *AB match-rho* **by** *auto*
qed
qed

qed
qed
next
case (*RuleB α Bs*)
from *RuleB*(*2*) **have** *lin*:*linear-term* (*lhs α*)
using *left-lin left-linear-trs-def* **by** *fastforce*
from *RuleB*(*2*) **obtain** *f′ ts* **where** *lhs*:*lhs α* = *Fun f′ ts*
using *no-var-lhs* **by** *fastforce*
show *?thesis* **proof**(*cases A ⊔ B*)
case *None*
with *subterm*(*4*) *RuleB* **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC x*)
with *subterm*(*4*) *RuleB FunA* **show** *?thesis*
by (*metis None join-sym join-opt.simps*(*2*) *join-var-rule*)
next
case (*FunC Cs h*)
from *None* **show** *?thesis* **proof**(*cases B ⊔ C*)
case (*Some BC*)
obtain *BCs τ* **where** *τ*:*match C* (*to-pterm* (*lhs α*)) = *Some τ* **and** *BC*:*BC* = *Prule α BCs*
and *l-BCs*:*length BCs* = *length Bs* **and** *args-BC*:∀ *i* < *length Bs*. *Bs*!*i* ⊔ *τ* (*var-rule α* ! *i*) = *Some* (*BCs* ! *i*)
using *join-rule-fun*[*OF Some*[*unfolded RuleB FunC*] *subterm*(*3*)[*unfolded RuleB*]] *FunC*(*1*) **by** *blast*
with *None Some FunA* **show** *?thesis* **proof**(*cases match A* (*to-pterm* (*lhs α*)))
case (*Some σ*)
from *None* **obtain** *i* **where** *i*:*i* < *length* (*var-rule α*) *map2* (⊔) (*map σ* (*var-rule α*)) *Bs* ! *i* = *None*
unfolding *FunA RuleB join.simps Some*[*unfolded FunA*] *option.case*
by (*smt* (*verit*, *ccfv-threshold*) *length-map length-zip list-all-length min-less-iff-conj option.case-eq-if option.distinct*(*1*) *those-not-none-xs*)
let *?x*=*var-rule α* ! *i*
have *subt*:*A* ⊳ *σ ?x* **using** *lhs i Some*
by (*smt* (*verit*, *ccfv-SIG*) *lin linear-term-var-vars-term-list match-matches nth-mem set-vars-term-list subst-image-subterm to-pterm.simps*(*2*) *vars-to-pterm*)

have *wf-τ-x*:*τ ?x* ∈ *wf-pterm R*
using *subterm*(*4*) *τ i*(*1*) **by** (*metis match-well-def vars-to-pterm*)
have *IH*:*join-opt* (*σ ?x*) (*Bs* ! *i* ⊔ *τ ?x*) = *join-opt* (*τ ?x*) (*σ ?x* ⊔ *Bs* ! *i*)
using *subterm*(*1*) *i wf-τ-x* **by** (*metis RuleB*(*3*) *RuleB*(*4*) *Some match-well-def nth-mem subt subterm.prems*(*1*) *vars-to-pterm*)
have (*Bs* ! *i* ⊔ *τ ?x*) = *Some* (*BCs* ! *i*)
using *args-BC i RuleB*(*3*) **by** *auto*
with *IH i* **have** (*σ ?x*) ⊔ (*BCs* ! *i*) = *None*
by (*simp add*: *RuleB*(*3*))
then have (*map2* (⊔) (*map σ* (*var-rule α*)) *BCs*) ! *i* = *None*

using *l-BCs i* **by** (*simp add*: *RuleB*(*3*))
then have *those* (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *BCs*) = *None*
using *l-BCs i* **by** (*metis* (*no-types*, *opaque-lifting*) *RuleB*(*3*) *length-map length-zip min.idem nth-mem option.exhaust those-not-none-x*)
then have $A \sqcup BC$ = *None*
using *Some i* **unfolding** *BC FunA join.simps* **by** *simp*
then show *?thesis*
unfolding *None* ‹$B \sqcup C$ = *Some BC*› **by** *auto*
qed *simp*
qed *simp*
next
case (*RuleC* β *Cs*)
from *None* **show** *?thesis* **proof**(*cases* $B \sqcup C$)
case (*Some BC*)
then have $\alpha\beta$:$\alpha = \beta$ **and** *l-B-C*:*length Bs* = *length Cs*
using *join-rule-rule*[*OF Some*[*unfolded RuleB RuleC*] *subterm*(*3*,*4*)[*unfolded RuleB RuleC*]] **by** *simp+*
from *Some* **obtain** *BCs* **where** *BC*:*BC* = *Prule* α *BCs* **and** *l-BC-B*:*length BCs* = *length Bs*
and *args-BC*:($\forall$ *i*<*length Cs*. *Bs*!*i* $\sqcup$ *Cs*!*i* = *Some* (*BCs* ! *i*))
using *join-rule-rule*[*OF Some*[*unfolded RuleB RuleC*] *subterm*(*3*,*4*)[*unfolded RuleB RuleC*]] **by** *force*
from *Some FunA RuleB BC* **show** *?thesis* **proof**(*cases match A* (*to-pterm* (*lhs* α)))
case (*Some* σ)
from *None* **obtain** *i* **where** *i*:*i* < *length* (*var-rule* α) *map2* ($\sqcup$) (*map* σ (*var-rule* α)) *Bs* ! *i* = *None*
unfolding *FunA RuleB join.simps Some*[*unfolded FunA*] *option.case*
by (*smt* (*verit*, *ccfv-threshold*) *length-map length-zip list-all-length min-less-iff-conj option.case-eq-if option.distinct*(*1*) *those-not-none-xs*)
let *?x*=*var-rule* α ! *i*
have *subt*:$A \rhd \sigma$ *?x* **using** *lhs i Some*
by (*smt* (*verit*, *ccfv-SIG*) *lin linear-term-var-vars-term-list match-matches nth-mem set-vars-term-list subst-image-subterm to-pterm.simps*(*2*) *vars-to-pterm*)

have *IH*:*join-opt* (σ *?x*) (*Bs* ! *i* $\sqcup$ *Cs* ! *i*) = *join-opt* (*Cs* ! *i*) (σ *?x* $\sqcup$ *Bs* ! *i*)
using *subterm*(*1*) *i RuleC* **by** (*metis RuleB*(*3*) *RuleB*(*4*) *Some* $\alpha\beta$ *match-well-def nth-mem subt subterm.prems*(*1*) *vars-to-pterm*)
from *IH i* **have** (σ *?x*) $\sqcup$ (*BCs* ! *i*) = *None*
using *RuleB*(*3*) *args-BC l-B-C* **by** *auto*
then have (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *BCs*) ! *i* = *None*
using *RuleB*(*3*) *i*(*1*) *l-BC-B* **by** *force*
then have *those* (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *BCs*) = *None*
by (*metis* (*no-types*, *opaque-lifting*) *RuleB*(*3*) *i*(*1*) *l-BC-B length-map length-zip min.idem not-Some-eq nth-mem those-not-none-x*)
then have $A \sqcup BC$ = *None*
using *Some i* **unfolding** *BC FunA join.simps* **by** *simp*
then show *?thesis*

unfolding *None* ‹$B \sqcup C$ = *Some BC*› **by** *auto*
qed *simp*
qed *simp*
qed
next
case (*Some AB*)
then obtain σ *ABs* **where** *sigma*:*match A* (*to-pterm* (*lhs* α)) = *Some* σ
and *AB*:*AB* = *Prule* α *ABs* **and** *l-ABs*:*length ABs* = *length Bs*
and *args-AB*:($\forall$ *i*<*length Bs*. σ (*var-rule* α ! *i*) $\sqcup$ *Bs* ! *i* = *Some* (*ABs* ! *i*))
unfolding *FunA RuleB* **using** *join-sym join-rule-fun subterm*(*2*,*3*)[*unfolded FunA RuleB*] *RuleB*(*3*) **by** (*smt* (*verit*, *del-insts*))
{**fix** *i* **assume** *i*:*i* < *length Bs*
with *sigma RuleB*(*3*) **have** (*map* σ (*var-rule* α))!*i* $\lhd$ *A*
using *lhs* **by** (*smt* (*verit*, *ccfv-threshold*) *lin linear-term-var-vars-term-list match-matches nth-map nth-mem set-vars-term-list subst-image-subterm to-pterm.simps*(*2*) *vars-to-pterm*)
}**note** *A-sub*=*this*
from *subterm*(*4*) **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC x*)
have *match* (*Var x*) (*to-pterm* (*lhs* α)) = *None*
unfolding *lhs to-pterm.simps* **using** *match-matches not-None-eq* **by** *fastforce*
then show *?thesis*
unfolding *Some* **unfolding** *RuleB VarC AB* **by** *simp*
next
case (*FunC Cs g*)
show *?thesis* **proof**(*cases match C* (*to-pterm* (*lhs* α)))
case *None*
then have $B \sqcup C$ = *None*
unfolding *RuleB FunC* **by** *simp*
moreover from *None* **have** $AB \sqcup C$ = *None*
unfolding *AB FunC* **by** *simp*
ultimately show *?thesis*
unfolding *Some* **by** (*simp add*: *join-sym*)
next
case (*Some* τ)
{**fix** *i* **assume** *i*:*i* < *length Bs*
have (*map* σ (*var-rule* α))!*i* $\in$ *wf-pterm R*
using *i match-well-def*[*OF subterm*(*2*) *sigma*] *RuleB*(*3*) **by** (*simp add*: *vars-to-pterm*)
moreover have (*map* τ (*var-rule* α))!*i* $\in$ *wf-pterm R*
using *i match-well-def*[*OF subterm*(*4*) *Some*] *RuleB*(*3*) **by** (*simp add*: *vars-to-pterm*)
ultimately have *join-opt* (*map* σ (*var-rule* α) ! *i*) (*Bs*!*i* $\sqcup$ *map* τ (*var-rule* α) ! *i*) = *ABs* ! *i* $\sqcup$ *map* τ (*var-rule* α) ! *i*
using *subterm*(*1*) *RuleB*(*3*,*4*) *i args-AB A-sub join-sym* **by** *fastforce*
}**note** *IH*=*this*
show *?thesis* **proof**(*cases those* (*map2* ($\sqcup$) *Bs* (*map* τ (*var-rule* α))))

case *None*
then obtain i **where** i:i < *length Bs Bs*!i ⊔ (*map* τ (*var-rule* α))!i = *None*
using *those-not-none-xs* **by** (*smt* (*verit*) *length-map length-zip list-all-length map-nth-eq-conv min-less-iff-conj nth-zip old.prod.case*)
with *IH* **have** *map* τ (*var-rule* α) ! i ⊔ *ABs* ! i = *None*
using *join-sym* **by** (*metis join-opt.simps*(*2*))
with i *RuleB*(*3*) *l-ABs* **have** i < *length* (*map2* (⊔) (*map* τ (*var-rule* α)) *ABs*) (*map2* (⊔) (*map* τ (*var-rule* α)) *ABs*) ! i = *None* **by** *simp-all*
then have *those* (*map2* (⊔) (*map* τ (*var-rule* α)) *ABs*) = *None*
by (*metis nth-mem option.exhaust those-not-none-x*)
with *None* **show** *?thesis*
unfolding ‹*A* ⊔ *B* = *Some AB*› *AB* **unfolding** *RuleB FunC join-opt.simps join.simps Some*[*unfolded FunC*] *option.case None* **by** *simp*
next
case (*Some BCs*)
then have *BC*:*B* ⊔ *C* = *Some* (*Prule* α *BCs*)
unfolding *RuleB FunC join.simps* ‹*match C* (*to-pterm* (*lhs* α)) = *Some* τ›[*unfolded FunC*] **by** *simp*
from *Some* **have** *l-BCs*:*length BCs* = *length Bs*
using *RuleB*(*3*) *length-those* **by** *fastforce*
{**fix** i **assume** i < *length Bs*
with *Some IH* **have** (*map* σ (*var-rule* α)) ! i ⊔ *BCs*!i = (*map* τ (*var-rule* α)) ! i ⊔ *ABs* ! i
using *RuleB*(*3*) *those-some2 join-sym* **by** *fastforce*
}
then have *map2* (⊔) (*map* σ (*var-rule* α)) *BCs* = *map2* (⊔) (*map* τ (*var-rule* α)) *ABs*
using *l-BCs l-ABs* **by** (*simp add*: *RuleB*(*3*) *map-eq-conv*′)
then show *?thesis* **unfolding** *BC* ‹*A* ⊔ *B* = *Some AB*› *AB* **unfolding** *FunA RuleB FunC AB join-opt.simps join.simps sigma*[*unfolded FunA*]
‹*match C* (*to-pterm* (*lhs* α)) = *Some* τ›[*unfolded FunC*] *option.case* **by** *simp*
qed
qed
next
case (*RuleC* β *Cs*)
show *?thesis* **proof**(*cases* $\alpha = \beta$)
case *True*
with *RuleB*(*3*) *RuleC*(*3*) **have** *l-Bs-Cs*:*length Bs* = *length Cs* **by** *simp*
{**fix** i **assume** i:i < *length Bs*
have (*map* σ (*var-rule* α))!i ∈ *wf-pterm R*
using i *match-well-def*[*OF subterm*(*2*) *sigma*] *RuleB*(*3*) **by** (*simp add*: *vars-to-pterm*)
then have *join-opt* (*map* σ (*var-rule* α) ! i) (*Bs*!i ⊔ *Cs* ! i) = *Cs* ! i ⊔ *ABs* ! i
using *subterm*(*1*) *RuleB*(*3*,*4*) *RuleC*(*3*,*4*) i *args-AB A-sub True* **by** *simp*
}**note** *IH*=*this*

```
          show ?thesis proof(cases those (map2 (⊔) Bs Cs))
           case None
           then obtain i where i:i < length Bs Bs!i ⊔ Cs!i = None
               using those-not-none-xs by (smt (verit) length-map length-zip
list-all-length map-nth-eq-conv min-less-iff-conj nth-zip old.prod.case)
           with IH have Cs ! i ⊔ ABs ! i = None by force
           with i RuleB(3) l-ABs l-Bs-Cs have i < length (map2 (⊔) Cs ABs)
(map2 (⊔) Cs ABs) ! i = None by simp-all
           then have those (map2 (⊔) Cs ABs) = None
            by (metis nth-mem option.exhaust those-not-none-x)
          with None show ?thesis unfolding ‹A ⊔ B = Some AB› AB unfolding
RuleB RuleC by simp
          next
           case (Some BCs)
           then have BC:B ⊔ C = Some (Prule α BCs)
            unfolding RuleB RuleC True by simp
           from Some have l-BCs:length BCs = length Bs
            using l-Bs-Cs length-those by fastforce
           {fix i assume i < length Bs
             with Some IH have (map σ (var-rule α)) ! i ⊔ BCs!i = Cs ! i ⊔
ABs ! i
             using those-some2 l-Bs-Cs by fastforce
           }
           then have map2 (⊔) (map σ (var-rule α)) BCs = map2 (⊔) Cs ABs
               using l-Bs-Cs RuleB(3) l-ABs l-BCs by (simp add: RuleC(3)
map-eq-conv')
           then show ?thesis
             unfolding BC ‹A ⊔ B = Some AB› AB unfolding FunA RuleB
RuleC join-opt.simps join.simps sigma[unfolded FunA] unfolding True by simp
          qed
        next
          case False
          then show ?thesis
          unfolding ‹A ⊔ B = Some AB› unfolding RuleB RuleC AB join.simps
by simp
        qed
      qed
    qed
  qed
 next
  case (RuleA α As)
  from RuleA(2) have lin:linear-term (lhs α)
    using left-lin left-linear-trs-def by fastforce
  from RuleA(2) obtain f' ts where lhs:lhs α = Fun f' ts
    using no-var-lhs by fastforce
 from subterm(3,2) show ?thesis proof(cases B rule:wf-pterm.cases[case-names
VarB FunB RuleB])
    case (VarB x)
    have match (Var x) (to-pterm (lhs α)) = None
```

unfolding *lhs* **using** *match-matches not-Some-eq* **by** *fastforce*
then show *?thesis* **unfolding** *RuleA VarB*
by (*metis join-sym RuleA*(*2*) *join.simps*(*1*) *join.simps*(*9*) *join-opt.simps*(*1*) *join-opt.simps*(*2*)
left-lin-no-var-lhs.join-var-rule left-lin-no-var-lhs-axioms subterm.prems(*3*) *wf-pterm.simps*)
next
case (*FunB Bs f*)
show *?thesis* **proof**(*cases A ⊔ B*)
case *None*
with *subterm*(*4*) *FunB* **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC x*)
with *subterm*(*4*) *FunB RuleA None* **show** *?thesis*
by *auto*
next
case (*FunC Cs h*)
from *None* **show** *?thesis* **proof**(*cases B ⊔ C*)
case (*Some BC*)
have *fh*:*f* = *h* **and** *l-B-C*:*length Bs* = *length Cs*
using *join-fun-fun*[*OF Some*[*unfolded FunB FunC*]] **by** *simp+*
obtain *BCs* **where** *BC*:*BC* = *Pfun f BCs* **and** *l-BC-B*:*length BCs* = *length Bs*
and *args-BC*:(∀ *i*<*length Bs. Bs*!*i* ⊔ *Cs*!*i* = *Some* (*BCs* ! *i*))
using *join-fun-fun*[*OF Some*[*unfolded FunB FunC*]] **by** *blast*
show *?thesis* **proof**(*cases match B* (*to-pterm* (*lhs α*)))
case *None*
then have ¬ *matches BC* (*to-pterm* (*lhs α*))
using *join-pterm-linear-subst* ‹*B* ⊔ *C* = *Some BC*› *lin* **by** (*metis match-complete′ matches-iff option.simps*(*3*))
then have *A* ⊔ *BC* = *None* **unfolding** *RuleA BC*
by (*smt* (*verit*) *join.simps*(*5*) *match-matches matches-iff not-Some-eq option.simps*(*4*))
then show *?thesis*
unfolding ‹*A* ⊔ *B* = *None*› ‹*B* ⊔ *C* = *Some BC*› **by** *simp*
next
case (*Some σ*)
with *None* **have** *those* (*map2* (⊔) *As* (*map σ* (*var-rule α*))) = *None*
unfolding *RuleA FunB* **using** *not-None-eq* **by** *fastforce*
then obtain *i* **where** *i*:*i* < *length* (*var-rule α*) *map2* (⊔) *As* (*map σ* (*var-rule α*)) ! *i* = *None*
by (*smt* (*verit, best*) *RuleA*(*3*) *length-map length-zip list-all-length min.idem those-not-none-xs*)
let *?x*=*var-rule α* ! *i*
from *i* **have** *none-at-i*:*As* ! *i* ⊔ *σ ?x* = *None*
using *RuleA*(*3*) **by** *simp*
show *?thesis* **proof**(*cases match C* (*to-pterm* (*lhs α*)))
case *None*
then have ¬ *matches BC* (*to-pterm* (*lhs α*))

using *join-pterm-linear-subst* ‹*B* ⊔ *C* = *Some BC*› *lin* **by** (*metis match-complete′ matches-iff option.simps*(*3*))
then have *A* ⊔ *BC* = *None* **unfolding** *RuleA BC*
by (*smt* (*verit*) *join.simps*(*5*) *match-matches matches-iff not-Some-eq option.simps*(*4*))
then show *?thesis*
unfolding ‹*A* ⊔ *B* = *None*› ‹*B* ⊔ *C* = *Some BC*› **by** *simp*
next
case (*Some* τ)
then obtain ϱ **where** *rho*:(∀ *x*∈*vars-term* (*to-pterm* (*lhs* α)). σ *x* ⊔ τ *x* = *Some* (ϱ *x*))
and *BC-rho*:*BC* = (*to-pterm* (*lhs* α)) · ϱ **and** *match-rho*:*match BC* (*to-pterm* (*lhs* α)) = *Some* ϱ
using *join-pterm-subst-Some match-matches* ‹*match B* (*to-pterm* (*lhs* α)) = *Some* σ› ‹*B* ⊔ *C* = *Some BC*› **by** *blast*
have σ *?x* ∈ *wf-pterm R*
using *i*(*1*) ‹*match B* (*to-pterm* (*lhs* α)) = *Some* σ› *subterm*(*3*) **by** (*metis match-well-def vars-to-pterm*)
moreover have τ *?x* ∈ *wf-pterm R*
using *i*(*1*) *Some subterm*(*4*) **by** (*metis match-well-def vars-to-pterm*)

ultimately have *IH*:*join-opt* (*As* ! *i*) (σ *?x* ⊔ τ *?x*) = *join-opt* (τ *?x*) (*As* ! *i* ⊔ σ *?x*)
using *subterm*(*1*) *i*(*1*) *RuleA*(*3*) **by** (*metis RuleA*(*1*) *RuleA*(*4*) *nth-mem supt.arg*)
then have (*As* ! *i*) ⊔ (ϱ *?x*) = *None*
using *none-at-i rho* **by** (*metis i*(*1*) *join-opt.simps*(*1*) *join-opt.simps*(*2*) *lin linear-term-var-vars-term-list nth-mem set-vars-term-list vars-to-pterm*)
then have (*map2* (⊔) *As* (*map* ϱ (*var-rule* α))) ! *i* = *None*
using *RuleA*(*3*) *i*(*1*) **by** *auto*
then have *those* (*map2* (⊔) *As* (*map* ϱ (*var-rule* α))) = *None*
by (*metis* (*no-types*, *opaque-lifting*) *RuleA*(*3*) *i*(*1*) *length-map length-zip min.idem not-Some-eq nth-mem those-not-none-x*)
then have *A* ⊔ *BC* = *None*
using *BC RuleA*(*1*) *match-rho* **by** *force*
then show *?thesis*
unfolding ‹*A* ⊔ *B* = *None*› ‹*B* ⊔ *C* = *Some BC*› **by** *simp*
qed
qed
qed *simp*
next
case (*RuleC* β *Cs*)
from *None* **show** *?thesis* **proof**(*cases B* ⊔ *C*)
case (*Some BC*)
obtain *BCs* τ **where** τ:*match B* (*to-pterm* (*lhs* β)) = *Some* τ **and** *BC*:*BC* = *Prule* β *BCs*
and *l-BCs*:*length BCs* = *length Cs* **and** *args-BC*:∀ *i* < *length Cs*. τ (*var-rule* β ! *i*) ⊔ *Cs*!*i* = *Some* (*BCs* ! *i*)
using *join-rule-fun Some*[*unfolded RuleC FunB*] *subterm*(*4*)[*unfolded*

RuleC] *FunB*(*1*) **by** (*metis Residual-Join-Deletion.join-sym*)
show *?thesis* **proof**(*cases match B* (*to-pterm* (*lhs* α)))
case *None*
with ‹*A* $\sqcup$ *B* = *None*› *Some BC RuleA*(*1*) τ **show** *?thesis* **by** *fastforce*
next
case (*Some* σ)
from *None* **obtain** *i* **where** *i*:*i* < *length* (*var-rule* α) *map2* ($\sqcup$) *As* (*map* σ (*var-rule* α)) ! *i* = *None*
unfolding *FunB RuleA join.simps Some*[*unfolded FunB*] *option.case*
by (*smt* (*verit*, *ccfv-threshold*) *length-map length-zip list-all-length min-less-iff-conj option.case-eq-if option.distinct*(*1*) *those-not-none-xs*)
let *?x*=*var-rule* α ! *i*
have *wf-*σ*-x*:σ *?x* $\in$ *wf-pterm R*
using *subterm*(*3*) *Some i*(*1*) **by** (*metis match-well-def vars-to-pterm*)

from *BC None* ‹*B* $\sqcup$ *C* = *Some BC*› *RuleA* **show** *?thesis* **proof**(*cases* $\alpha = \beta$)
case *True*
then have $\sigma = \tau$
using *Some* τ **by** *auto*
have *IH*:*join-opt* (*As* ! *i*) (τ *?x* $\sqcup$ *Cs* ! *i*) = *join-opt* (*Cs* ! *i*) (*As* ! *i* $\sqcup$ τ *?x*)
using *subterm*(*1*) *i wf-*σ*-x args-BC* **by** (*metis RuleA*(*1*) *RuleA*(*3*) *RuleA*(*4*) *RuleC*(*3*) *RuleC*(*4*) *True* ‹$\sigma = \tau$› *nth-mem supt.arg*)
have τ *?x* $\sqcup$ *Cs* ! *i* = *Some* (*BCs* ! *i*)
using *args-BC i RuleC*(*3*) *True* **by** *force*
with *IH i* **have** (*As* ! *i*) $\sqcup$ (*BCs* ! *i*) = *None*
by (*simp add*: *RuleA*(*3*) ‹$\sigma = \tau$›)
then have (*map2* ($\sqcup$) *As BCs*) ! *i* = *None*
using *l-BCs i* **by** (*simp add*: *RuleA*(*3*) *RuleC*(*3*) *True*)
then have *those* (*map2* ($\sqcup$) *As BCs*) = *None*
using *l-BCs i those-not-none-x* **by** (*metis RuleA*(*3*) *RuleC*(*3*) *True length-map length-zip min.idem nth-mem option.exhaust*)
then have *A* $\sqcup$ *BC* = *None*
by (*simp add*: *BC RuleA*(*1*))
then show *?thesis*
unfolding *None* ‹*B* $\sqcup$ *C* = *Some BC*› **by** *auto*
qed *simp*
qed
qed *simp*
qed
next
case (*Some AB*)
obtain σ *ABs* **where** *sigma*:*match B* (*to-pterm* (*lhs* α)) = *Some* σ
and *AB*:*AB* = *Prule* α *ABs* **and** *l-ABs*:*length ABs* = *length As*
and *args-AB*:($\forall$ *i*<*length As*. *As*!*i* $\sqcup$ σ (*var-rule* α ! *i*) = *Some* (*ABs* ! *i*))
using *join-sym join-rule-fun*[*OF Some*[*unfolded FunB RuleA*]] **using** *FunB*(*1*) *RuleA*(*1*) *subterm.prems*(*1*) **by** *blast*
from *subterm*(*4*) *FunB*(*1*) **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names*

VarC FunC RuleC])
case (*VarC x*)
have *match* (*Var x*) (*to-pterm* (*lhs* α)) = *None*
unfolding *lhs* **using** *match-matches not-Some-eq* **by** *fastforce*
then show *?thesis* **unfolding** *Some* **unfolding** *RuleA FunB AB VarC* **by** *simp*
next
case (*FunC Cs g*)
show *?thesis* **proof**(*cases* $f = g \wedge$ *length Bs* = *length Cs*)
case *True*
show *?thesis* **proof**(*cases match C* (*to-pterm* (*lhs* α)))
case *None*
then have $*$:$C \sqcup AB$ = *None*
unfolding *AB FunC* **by** *simp*
with *Some* **show** *?thesis* **proof**(*cases* $B \sqcup C$)
case (*Some BC*)
with *None* **have** *match BC* (*to-pterm* (*lhs* α)) = *None*
by (*metis* (*no-types*, *lifting*) *domD domIff join-pterm-linear-subst lin match-complete′ match-lhs-subst option.simps*(*3*))
moreover obtain *BCs* **where** *BC* = *Pfun f BCs*
by (*metis FunB*(*1*) *FunC*(*1*) *Some join-fun-fun*)
ultimately show *?thesis*
using $*$ **unfolding** ‹$A \sqcup B$ = *Some AB*› *AB* **unfolding** *RuleA Some* **unfolding** *FunC* **by** (*simp add*: *join-sym*)
qed *simp*
next
case (*Some* τ)
{**fix** *i* **assume** *i*:$i <$ *length As*
have (*map* σ (*var-rule* α)) ! $i \in$ *wf-pterm R*
using *i match-well-def*[*OF subterm*(*3*) *sigma*] *RuleA*(*3*) **by** (*simp add*: *vars-to-pterm*)
moreover have (*map* τ (*var-rule* α)) ! $i \in$ *wf-pterm R*
using *i match-well-def*[*OF subterm*(*4*) *Some*] *RuleA*(*3*) **by** (*simp add*: *vars-to-pterm*)
ultimately have *join-opt* (*As* ! *i*) ((*map* σ (*var-rule* α) ! *i*) $\sqcup$ (*map* τ (*var-rule* α) ! *i*)) = (*map* τ (*var-rule* α) ! *i*) $\sqcup$ *ABs* ! *i*
using *subterm*(*1*) *RuleA*(*1*,*3*,*4*) *i args-AB True* **by** (*metis* (*no-types*, *lifting*) *join-opt.simps*(*1*) *nth-map nth-mem supt.arg*)
}**note** *IH*=*this*
show *?thesis* **proof**(*cases* $B \sqcup C$)
case *None*
with *sigma Some* **obtain** *x* **where** $x \in$ *vars-term* (*lhs* α) **and** $\sigma\ x \sqcup \tau\ x$ = *None*
using *join-pterm-subst-None* **by** (*metis lhs-subst-trivial match-lhs-subst option.sel set-vars-term-list vars-to-pterm*)
then obtain *i* **where** *i*:$i <$ *length* (*var-rule* α) (*map* σ (*var-rule* α) ! *i*) $\sqcup$ (*map* τ (*var-rule* α) ! *i*) = *None*
by (*metis* (*no-types*, *opaque-lifting*) *in-set-idx lin linear-term-var-vars-term-list nth-map set-vars-term-list*)

with *IH* **have** *map2* $(\sqcup)$ (*map* τ (*var-rule* α)) *ABs* ! *i* = *None*
using *RuleA*(*3*) *l-ABs* **by** *fastforce*
with *i*(*1*) **have** *those* (*map2* $(\sqcup)$ (*map* τ (*var-rule* α)) *ABs*) = *None*
using *those-not-none-x* **by** (*metis* (*no-types*, *opaque-lifting*) *RuleA*(*3*) *l-ABs length-map length-zip min.idem nth-mem option.exhaust*)
with *Some* **show** *?thesis* **unfolding** ‹*A* $\sqcup$ *B* = *Some AB*› *AB None* **unfolding** *FunC* **by** *simp*
next
case (*Some BC*)
from *sigma Some* **obtain** ϱ **where** *rho*:($\forall$ *x*$\in$*vars-term* (*to-pterm* (*lhs* α)). σ *x* $\sqcup$ τ *x* = *Some* (ϱ *x*))
and *BC-rho*:*BC* = (*to-pterm* (*lhs* α)) $\cdot$ ϱ **and** *match-rho*:*match BC* (*to-pterm* (*lhs* α)) = *Some* ϱ
using *join-pterm-subst-Some match-matches* ‹*match C* (*to-pterm* (*lhs* α)) = *Some* τ› **by** *blast*
{**fix** *i* **assume** *i*:*i* < *length As*
with *rho* **have** (*map* σ (*var-rule* α)) ! *i* $\sqcup$ (*map* τ (*var-rule* α)) ! *i* = *Some* ((*map* ϱ (*var-rule* α)) ! *i*)
by (*metis* (*no-types*, *lifting*) *RuleA*(*3*) *lin linear-term-var-vars-term-list nth-map nth-mem set-vars-term-list vars-to-pterm*)
with *i IH* **have** *As* ! *i* $\sqcup$ (*map* ϱ (*var-rule* α)) ! *i* = (*map* τ (*var-rule* α)) ! *i* $\sqcup$ *ABs* ! *i*
by *force*
}
then have *map2* $(\sqcup)$ *As* (*map* ϱ (*var-rule* α)) = *map2* $(\sqcup)$ (*map* τ (*var-rule* α)) *ABs*
by (*simp add*: *RuleA*(*3*) *l-ABs map-equality-iff*)
with *match-rho* ‹*match C* (*to-pterm* (*lhs* α)) = *Some* τ›
show *?thesis* **unfolding** ‹*A* $\sqcup$ *B* = *Some AB*› *AB Some*
unfolding *RuleA BC-rho join-opt.simps to-pterm.simps FunC* **by** (*simp add*: *lhs*)
qed
qed
next
case *False*
then consider (*fg*) *f* $\neq$ *g* | (*length*) *length Bs* $\neq$ *length Cs* **by** *fastforce*
then show *?thesis* **proof**(*cases*)
case *fg*
from *sigma* **have** $f' = f$
unfolding *FunB lhs to-pterm.simps* **using** *match-matches* **by** *fastforce*

with *fg* **have** *match C* (*to-pterm* (*lhs* α)) = *None*
unfolding *lhs FunC* **using** *domIff match-matches* **by** *fastforce*
with *fg* **show** *?thesis* **unfolding** ‹*A* $\sqcup$ *B* = *Some AB*› **unfolding** *RuleA FunB FunC AB* **by** *simp*
next
case *length*
from *sigma* **have** *length ts* = *length Bs*
using *FunB*(*1*) *lhs match-matches* **by** *fastforce*

then have *match C* (*to-pterm* (*lhs* α)) = *None*
unfolding *FunC lhs* **using** *length*
by (*smt* (*verit*, *del-insts*) *eval-term.simps*(*2*) *length-map match-matches option.exhaust term.inject*(*2*) *to-pterm.simps*(*2*))
with *length* **show** *?thesis* **unfolding** ‹*A* ⊔ *B* = *Some AB*› **unfolding** *RuleA FunB FunC AB* **by** *simp*
qed
qed
next
case (*RuleC* β *Cs*)
show *?thesis* **proof**(*cases* α = β)
case *True*
{**fix** *i* **assume** *i*:*i* < *length As*
have (*map* σ (*var-rule* α))!*i* ∈ *wf-pterm R*
using *i match-well-def*[*OF subterm*(*3*) *sigma*] *RuleA*(*3*) **by** (*simp add*: *vars-to-pterm*)
then have *join-opt* (*As*!*i*) ((*map* σ (*var-rule* α) ! *i*) ⊔ *Cs* ! *i*) = *Cs* ! *i* ⊔ *ABs* ! *i*
using *subterm*(*1*) *RuleA*(*1*,*3*,*4*) *RuleC i args-AB True* **by** (*metis* (*no-types*, *lifting*) *join-opt.simps*(*1*) *nth-map nth-mem supt.arg*)
}**note** *IH*=*this*
show *?thesis* **proof**(*cases those* (*map2* (⊔) (*map* σ (*var-rule* α)) *Cs*))
case *None*
with *sigma* **have** *BC*:*B* ⊔ *C* = *None*
unfolding *FunB RuleC True* **by** *simp*
from *None* **obtain** *i* **where** *i*:*i* < *length* (*map2* (⊔) (*map* σ (*var-rule* α)) *Cs*) **and** (*map2* (⊔) (*map* σ (*var-rule* α)) *Cs*) ! *i* = *None*
using *list-all-length those-not-none-xs* **by** *blast*
with *IH* **have** (*map2* (⊔) *Cs ABs*)!*i* = *None*
using *RuleA*(*3*) *l-ABs RuleC*(*3*) **by** *fastforce*
with *i*(*1*) **have** *those* (*map2* (⊔) *Cs ABs*) = *None*
using *l-ABs RuleA*(*3*) *RuleC*(*3*) *those-not-none-x* **unfolding** *True*
by (*metis length-map length-zip not-Some-eq nth-mem*)
then show *?thesis* **unfolding** *BC* ‹*A* ⊔ *B* = *Some AB*› **unfolding** *AB RuleC True* **by** *simp*
next
case (*Some BCs*)
with *sigma* **have** *BC*:*B* ⊔ *C* = *Some* (*Prule* α *BCs*)
unfolding *FunB RuleC True* **by** *simp*
{**fix** *i* **assume** *i*:*i* < *length As*
with *Some* **have** (*map* σ (*var-rule* α)) ! *i* ⊔ *Cs* ! *i* = *Some* (*BCs* ! *i*)
using *RuleA*(*3*) *RuleC*(*3*) *True those-some2* **by** *fastforce*
with *i IH* **have** *As* ! *i* ⊔ *BCs* ! *i* = *Cs* ! *i* ⊔ *ABs* ! *i*
by *force*
}
moreover have *length Cs* = *length BCs*
using *RuleC*(*3*) *Some True length-those* **by** *fastforce*
ultimately have *map2* (⊔) *As BCs* = *map2* (⊔) *Cs ABs*
using *RuleA*(*3*) *l-ABs map-equality-iff*

by (*smt* (*verit*, *ccfv-threshold*) *RuleC*(*3*) *Some True length-map length-those length-zip nth-zip old.prod.case*)
then show *?thesis* **unfolding** ‹$A \sqcup B = Some\ AB$› *BC AB* **unfolding** *RuleA RuleC True* **by** *simp*
qed
next
case *False*
show *?thesis* **proof**(*cases* $B \sqcup C$)
case *None*
show *?thesis* **unfolding** *None* ‹$A \sqcup B = Some\ AB$› **unfolding** *AB RuleC* **using** *False* **by** *simp*
next
case (*Some BC*)
then obtain *BCs* **where** $BC = Prule\ \beta\ BCs$
unfolding *RuleC FunB* **by** (*metis Residual-Join-Deletion.join-sym RuleC*(*1*) *join-rule-fun subterm.prems*(*3*))
with *False* **show** *?thesis* **unfolding** ‹$A \sqcup B = Some\ AB$› *Some AB* **unfolding** *RuleC RuleA* **by** *simp*
qed
qed
qed
qed
next
case (*RuleB* β *Bs*)
show *?thesis* **proof**(*cases* $A \sqcup B$)
case *None*
then show *?thesis* **proof**(*cases* $\alpha = \beta$)
case *True*
then have *l-As-Bs*:*length As* = *length Bs*
by (*simp add*: *RuleA*(*3*) *RuleB*(*3*))
with *None* **obtain** *i* **where** $i{:}i < length\ As\ As\ !\ i \sqcup Bs\ !\ i = None$
unfolding *True RuleA RuleB* **unfolding** *join.simps* **by** (*smt* (*verit*) *RuleB*(*3*) *length-map length-zip list-all-length*
map-nth-eq-conv min-less-iff-conj nth-zip old.prod.case option.case-eq-if option.distinct(*1*) *those-not-none-xs*)
from *subterm*(*4*) **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC x*)
show *?thesis* **unfolding** *RuleA RuleB VarC*
by (*metis None Residual-Join-Deletion.join-sym RuleA*(*1*) *RuleB*(*1*) *RuleB*(*2*) *join-opt.simps*(*2*) *left-lin-no-var-lhs.join-var-rule left-lin-no-var-lhs-axioms*)
next
case (*FunC Cs f*)
from *None* **show** *?thesis* **proof**(*cases* $B \sqcup C$)
case (*Some BC*)
obtain *BCs* τ **where** $\tau{:}match\ C\ (to\text{-}pterm\ (lhs\ \beta)) = Some\ \tau$ **and** $BC{:}BC = Prule\ \beta\ BCs$
and *l-BCs*:*length BCs* = *length Bs* **and** $args\text{-}BC{:}\forall\, i < length\ Bs.\ Bs\ !\ i \sqcup \tau\ (var\text{-}rule\ \beta\ !\ i) = Some\ (BCs\ !\ i)$

using *join-rule-fun Some*[*unfolded RuleB FunC*] *subterm*(*3*)[*unfolded RuleB*] *FunC*(*1*) **by** *metis*
let *?x=var-rule* β ! *i*
have *IH*:*join-opt* (*As* ! *i*) (*Bs* ! *i* $\sqcup$ τ *?x*) = *join-opt* (τ *?x*) (*As* ! *i* $\sqcup$ *Bs* ! *i*)
by (*metis RuleA*(*1*) *RuleA*(*3*) *RuleA*(*4*) *RuleB*(*3*) *RuleB*(*4*) *True* τ *i*(*1*) *match-well-def nth-mem subterm.hyps subterm.prems*(*3*) *supt.arg vars-to-pterm*)
with *i* **have** *join-opt* (*As* ! *i*) (*Bs* ! *i* $\sqcup$ τ *?x*) = *None*
by *simp*
then have *As* ! *i* $\sqcup$ *BCs* ! *i* = *None*
using *args-BC i*(*1*) *l-As-Bs* **by** *auto*
then have (*map2* ($\sqcup$) *As BCs*) ! *i* = *None*
using *i*(*1*) *l-As-Bs l-BCs* **by** *force*
then have *those* (*map2* ($\sqcup$) *As BCs*) = *None*
by (*metis i*(*1*) *l-As-Bs l-BCs length-map length-zip min.idem not-None-eq nth-mem those-not-none-x*)
then show *?thesis*
using *None RuleA*(*1*) *True BC Some* **by** *auto*
qed *simp*
next
case (*RuleC* γ *Cs*)
from *None* **show** *?thesis* **proof**(*cases B* $\sqcup$ *C*)
case (*Some BC*)
then have $\beta = \gamma$
unfolding *RuleB RuleC* **by** (*metis join.simps*(*3*) *option.distinct*(*1*))
from *Some* **obtain** *BCs* **where** *BC*:*BC* = *Prule* β *BCs* **and** *l-BCs*:*length BCs* = *length Bs* **and**
args-BC:$\forall$ *i* < *length Bs*. *Bs* ! *i* $\sqcup$ *Cs* ! *i* = *Some* (*BCs* ! *i*)
using *RuleB*(*1*) *RuleC*(*1*) *join-rule-rule subterm.prems*(*2*) *subterm.prems*(*3*) **by** *blast*
have *IH*:*join-opt* (*As* ! *i*) (*Bs* ! *i* $\sqcup$ *Cs* ! *i*) = *join-opt* (*Cs* ! *i*) (*As* ! *i* $\sqcup$ *Bs* ! *i*)
using *subterm*(*1*) **by** (*metis RuleA*(*1*) *RuleA*(*3*) *RuleA*(*4*) *RuleB*(*4*) *RuleC*(*3*) *RuleC*(*4*) *True* ‹$\beta = \gamma$› *i*(*1*) *l-As-Bs nth-mem supt.arg*)
then have *As* ! *i* $\sqcup$ *BCs* ! *i* = *None*
using *args-BC i*(*1*) *l-As-Bs i*(*2*) **by** *fastforce*
then have (*map2* ($\sqcup$) *As BCs*) ! *i* = *None*
using *i*(*1*) *l-As-Bs l-BCs* **by** *force*
then have *those* (*map2* ($\sqcup$) *As BCs*) = *None*
by (*metis i*(*1*) *l-As-Bs l-BCs length-map length-zip min.idem not-None-eq nth-mem those-not-none-x*)
then show *?thesis*
using *None RuleA*(*1*) *True BC Some* **by** *auto*
qed *simp*
qed
next
case *False*
then show *?thesis* **proof**(*cases C*)
case (*Var x*)

show *?thesis* **unfolding** *RuleA RuleB Var*
by (*metis None Residual-Join-Deletion.join-sym RuleA*(*1*) *RuleB*(*1*) *RuleB*(*2*) *join-opt.simps*(*2*) *join-var-rule*)
next
case (*Pfun f Cs*)
show *?thesis* **unfolding** *RuleA RuleB Pfun*
by (*metis* (*no-types*, *lifting*) *False RuleB*(*1*) *join.simps*(*3*) *join-opt.elims join-opt.simps*(*2*) *join-rule-fun subterm.prems*(*2*))
next
case (*Prule* γ *Cs*)
show *?thesis* **unfolding** *RuleA RuleB Prule*
by (*smt* (*verit*, *ccfv-threshold*) *False Prule RuleA*(*1*) *RuleB*(*1*) *join-opt.elims join-rule-rule join-wf-pterm subterm.prems*(*1*) *subterm.prems*(*2*) *subterm.prems*(*3*))
qed
qed
next
case (*Some AB*)
then have *alpha-beta*:$\beta = \alpha$
unfolding *RuleA RuleB* **by** (*metis join.simps*(*3*) *option.distinct*(*1*))
with *Some* **obtain** *ABs* **where** *AB*:*AB* = *Prule* α *ABs* **and** *l-AB-A*:*length ABs* = *length As*
and *args-AB*:($\forall$ *i*<*length Bs*. *As*!*i* $\sqcup$ *Bs*!*i* = *Some* (*ABs* ! *i*))
by (*smt* (*verit*, *ccfv-SIG*) *RuleA*(*1*) *RuleB*(*1*) *join-rule-rule subterm.prems*(*1*) *subterm.prems*(*2*))
from *subterm*(*4*) **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC x*)
have *match* (*Var x*) (*to-pterm* (*Fun f' ts*)) = *None*
by (*metis case-optionE match-matches option.disc-eq-case*(*2*) *subst-apply-eq-Var term.distinct*(*1*) *to-pterm.simps*(*2*))
then show *?thesis* **unfolding** *Some AB* **unfolding** *RuleB VarC join.simps alpha-beta* **by** (*simp add*: *lhs*)
next
case (*FunC Cs f*)
show *?thesis* **proof**(*cases match C* (*to-pterm* (*lhs* α)))
case *None*
then show *?thesis* **unfolding** *Some AB* **unfolding** *RuleB alpha-beta FunC* **by** *simp*
next
case (*Some* σ)
{**fix** *i* **assume** *i*:*i* < *length As*
have (*map* σ (*var-rule* α))!*i* $\in$ *wf-pterm R*
using *i match-well-def*[*OF subterm*(*4*) *Some*] *RuleA*(*3*) **by** (*simp add*: *vars-to-pterm*)
with *i* **have** *join-opt* (*As* ! *i*) (*Bs* ! *i* $\sqcup$ (*map* σ (*var-rule* α) ! *i*)) = *map* σ (*var-rule* α) ! *i* $\sqcup$ *ABs* ! *i*
using *subterm*(*1*) *RuleA RuleB* **by** (*metis alpha-beta args-AB join-opt.simps*(*1*) *nth-mem supt.arg*)

}**note** *IH=this*
show *?thesis* **proof**(*cases those* (*map2* (⊔) *Bs* (*map* σ (*var-rule* α))))
case *None*
from *None* **obtain** *i* **where** *i*:*i* < *length* (*map2* (⊔) *Bs* (*map* σ (*var-rule* α))) **and** (*map2* (⊔) *Bs* (*map* σ (*var-rule* α))) ! *i* = *None*
using *list-all-length those-not-none-xs* **by** *blast*
with *IH* **have** (*map2* (⊔) (*map* σ (*var-rule* α)) *ABs*)!*i* = *None*
using *RuleA*(*3*) *l-AB-A* **by** *fastforce*
with *i*(*1*) **have** *those* (*map2* (⊔) (*map* σ (*var-rule* α)) *ABs*) = *None*
using *l-AB-A RuleA*(*3*) *those-not-none-x* **by** (*metis RuleB*(*3*) *alpha-beta length-map length-zip nth-mem option.exhaust*)
with ‹*A* ⊔ *B* = *Some AB*› *Some None* **show** *?thesis* **unfolding** *AB RuleB FunC alpha-beta* **by** *fastforce*
next
case (*Some BCs*)
then have *BC*:*B* ⊔ *C* = *Some* (*Prule* α *BCs*)
unfolding *RuleB FunC alpha-beta* **using** ‹*match C* (*to-pterm* (*lhs* α)) = *Some* σ› **by** (*simp add*: *FunC*(*1*))
then have *l-BCs*:*length BCs* = *length As*
using *RuleA*(*3*) *RuleB*(*3*) *Some alpha-beta length-those* **by** *force*
{**fix** *i* **assume** *i* < *length As*
then have *As*!*i* ⊔ *BCs*!*i* = (*map* σ (*var-rule* α)) ! *i* ⊔ *ABs* ! *i*
using *IH Some l-BCs length-those those-some2* **by** *fastforce*
}
then have *map2* (⊔) *As BCs* = *map2* (⊔) (*map* σ (*var-rule* α)) *ABs*
by (*simp add*: *RuleA*(*3*) *l-AB-A l-BCs map-equality-iff*)
then show *?thesis* **using** ‹*match C* (*to-pterm* (*lhs* α)) = *Some* σ›
unfolding ‹*A* ⊔ *B* = *Some AB*› *AB BC* **unfolding** *RuleA FunC join-opt.simps join.simps* **by** *simp*
qed
qed
next
case (*RuleC* γ *Cs*)
show *?thesis* **proof**(*cases* $\alpha = \gamma$)
case *True*
{**fix** *i* **assume** *i*:*i* < *length As*
then have *join-opt* (*As* ! *i*) (*Bs* ! *i* ⊔ *Cs* ! *i*) = *Cs* ! *i* ⊔ *ABs* ! *i*
using *subterm*(*1*) *RuleA RuleB RuleC True alpha-beta args-AB* **by** (*metis join-opt.simps*(*1*) *nth-mem supt.arg*)
}**note** *IH=this*
show *?thesis* **proof**(*cases those* (*map2* (⊔) *Bs Cs*))
case *None*
then obtain *i* **where** *i*:*i* < *length* (*map2* (⊔) *Bs Cs*) (*map2* (⊔) *Bs Cs*)!*i* = *None*
using *list-all-length those-not-none-xs* **by** *blast*
with *IH* **have** *map2* (⊔) *Cs ABs* ! *i* = *None*
using *RuleA*(*3*) *RuleC*(*3*) *True l-AB-A* **by** *fastforce*
with *i*(*1*) **have** *those* (*map2* (⊔) *Cs ABs*) = *None*
by (*metis RuleA*(*3*) *RuleB*(*3*) *RuleC*(*3*) *True alpha-beta l-AB-A*

length-map length-zip not-Some-eq nth-mem those-not-none-x)
with *None* **show** *?thesis* **unfolding** *Some AB* **unfolding** *RuleC RuleB True alpha-beta* **by** *simp*
next
case (*Some BCs*)
then have *BC*:$B \sqcup C$ = *Some* (*Prule* α *BCs*)
by (*simp add*: *RuleB*(*1*) *RuleC*(*1*) *True alpha-beta*)
{**fix** *i* **assume** $i <$ *length As*
with *IH* **have** *As*!*i* $\sqcup$ *BCs*!*i* = *Cs*!*i* $\sqcup$ *ABs*!*i*
using *RuleA*(*3*) *RuleB*(*3*) *RuleC*(*3*) *Some True alpha-beta those-some2* **by** *fastforce*
}
moreover have *length BCs* = *length As*
using *RuleA*(*3*) *RuleB*(*3*) *RuleC*(*3*) *Some True alpha-beta length-those* **by** *force*
ultimately have *those* (*map2* ($\sqcup$) *As BCs*) = *those* (*map2* ($\sqcup$) *Cs ABs*)
by (*smt* (*verit*, *ccfv-SIG*) *RuleA*(*3*) *RuleB*(*3*) *RuleC*(*3*) *Some True alpha-beta l-AB-A length-map length-those length-zip map-equality-iff nth-zip old.prod.case*)
then show *?thesis* **unfolding** ‹$A \sqcup B$ = *Some AB*› *AB BC* **unfolding** *RuleA RuleC True alpha-beta* **by** *simp*
qed
next
case *False*
then show *?thesis* **unfolding** ‹$A \sqcup B$ = *Some AB*› *AB* **unfolding** *RuleA RuleB RuleC*
by (*simp add*: *alpha-beta*)
qed
qed
qed
qed
qed
qed

Preparation for well-definedness result for $\bigsqcup$.

lemma *join-triple-defined*:
assumes $A \in$ *wf-pterm R* $B \in$ *wf-pterm R* $C \in$ *wf-pterm R*
and $A \sqcup B \neq$ *None* $B \sqcup C \neq$ *None* $A \sqcup C \neq$ *None*
shows *join-opt A* ($B \sqcup C$) $\neq$ *None*
using *assms* **proof**(*induct A arbitrary*:*B C rule*:*subterm-induct*)
case (*subterm A*)
from *subterm*(*5*) **obtain** *AB* **where** *joinAB*:$A \sqcup B$ = *Some AB* **by** *blast*
from *subterm*(*6*) **obtain** *BC* **where** *joinBC*:$B \sqcup C$ = *Some BC* **by** *blast*
from *subterm*(*7*) **obtain** *AC* **where** *joinAC*:$A \sqcup C$ = *Some AC* **by** *blast*
from *subterm*(*2*) **show** *?case* **proof**(*cases A rule*:*wf-pterm.cases*[*case-names VarA FunA RuleA*])
case (*VarA x*)
from *subterm*(*3*,*5*) **show** *?thesis* **proof**(*cases B rule*:*wf-pterm.cases*[*case-names*

VarB FunB RuleB])
case (*VarB y*)
from *subterm*(*5*) **have** *x*:$x = y$ **unfolding** *VarA VarB* **by** (*meson join.simps*(*1*))

from *subterm*(*4*,*6*) **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC z*)
from *subterm*(*6*) **show** *?thesis* **unfolding** *VarA VarB x VarC*
by (*metis join.simps*(*1*) *join-opt.simps*(*1*))
next
case (*RuleC* α *Cs*)
from *subterm*(*5*−) **show** *?thesis* **unfolding** *VarA VarB RuleC x*
by (*metis Residual-Join-Deletion.join-sym RuleC*(*1*) *VarA join-opt.elims join-with-source option.sel source.simps*(*1*) *source-join subterm.prems*(*1*) *subterm.prems*(*3*) *to-pterm.simps*(*1*) *x*)
qed (*simp add*: *VarB*)
next
case (*RuleB* α *Bs*)
from *subterm*(*2*−) *VarA no-var-lhs RuleB* **show** *?thesis*
by (*metis join-sym join-opt.elims join-wf-pterm join-with-source source.simps*(*1*) *source-join to-pterm.simps*(*1*))
qed (*simp add*: *VarA*)
next
case (*FunA As f*)
from *subterm*(*3*,*5*) **show** *?thesis* **proof**(*cases B rule*:*wf-pterm.cases*[*case-names VarB FunB RuleB*])
case (*FunB Bs g*)
from *subterm*(*5*) **have** *fg*:$f = g$ **and** *l-A-B*:*length As* = *length Bs*
unfolding *FunA FunB* **by** (*meson join.simps*(*2*))+
obtain *ABs* **where** *AB*:*AB* = *Pfun f ABs* **and** *l-AB-A*:*length ABs* = *length As*
and *args-AB*:($\forall$ *i*<*length Bs*. *As*!*i* $\sqcup$ *Bs*!*i* = *Some* (*ABs* ! *i*))
using *join-fun-fun*[*OF joinAB*[*unfolded FunA FunB*]] **by** *fastforce*
from *subterm*(*4*,*6*) **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*FunC Cs h*)
from *subterm*(*6*) **have** *gh*:$g = h$ **and** *l-B-C*:*length Bs* = *length Cs*
unfolding *FunB FunC* **by** (*meson join.simps*(*2*))+
from *subterm*(*7*) **have** *fh*:$f = h$ **and** *l-A-C*:*length As* = *length Cs*
unfolding *FunA FunC* **by** (*meson join.simps*(*2*))+
obtain *BCs* **where** *BC*:*BC* = *Pfun g BCs* **and** *l-BC-B*:*length BCs* = *length Bs*
and *args-BC*:($\forall$ *i*<*length BCs*. *Bs*!*i* $\sqcup$ *Cs*!*i* = *Some* (*BCs* ! *i*))
using *join-fun-fun*[*OF joinBC*[*unfolded FunB FunC*]] **by** *fastforce*
obtain *ACs* **where** *AC*:*AC* = *Pfun h ACs* **and** *l-AC-C*:*length ACs* = *length Cs*
and *args-AC*:($\forall$ *i*<*length ACs*. *As*!*i* $\sqcup$ *Cs*!*i* = *Some* (*ACs* ! *i*))
using *join-fun-fun*[*OF joinAC*[*unfolded FunA FunC*]] **by** *fastforce*
have *those* (*map2* ($\sqcup$) *As BCs*) $\neq$ *None* **proof**−

{**fix** i **assume** i:i < *length* (*zip As BCs*)
from *FunA FunB FunC* i **have** *join-opt* (*As*!i) ((*Bs*!i) ⊔ (*Cs*!i)) ≠ *None*
using *subterm*(*1*) *l-A-B l-B-C l-AC-C* **by** (*smt* (*verit*, *ccfv-threshold*) *args-AB args-AC args-BC length-zip min-less-iff-conj nth-mem option.distinct*(*1*) *supt.arg*)
then have (*map2* (⊔) *As BCs*)!i ≠ *None*
using i *args-BC* **by** *simp*
}
then show *?thesis*
by (*simp add*: *list-all-length those-not-none-xs*)
qed
then show *?thesis*
unfolding *joinBC BC* **unfolding** *FunA fg gh join-opt.simps*
by (*simp add*: *l-A-B l-BC-B option.case-eq-if*)
next
case (*RuleC* α *Cs*)
from *joinBC subterm*(*4*) **obtain** σ *BCs* **where** *match-lhs-B*:*match B* (*to-pterm* (*lhs* α)) = *Some* σ
and *BC*:*BC* = *Prule* α *BCs* **and** *l-BC-C*:*length BCs* = *length Cs*
and *args-BC*:(∀ *i*<*length Cs*. *Cs* ! *i* ⊔ σ (*var-rule* α ! *i*) = *Some* (*BCs* ! *i*))
unfolding *FunB RuleC* **using** *join-rule-fun RuleC*(*1*,*2*,*3*) *join-sym* **by** *metis*
from *joinAC subterm*(*4*) **obtain** τ *ACs* **where** *match-lhs-A*:*match A* (*to-pterm* (*lhs* α)) = *Some* τ
and *AC*:*AC* = *Prule* α *ACs* **and** *l-AC-C*:*length ACs* = *length Cs*
and *args-AC*:(∀ *i*<*length Cs*. *Cs* ! *i* ⊔ τ (*var-rule* α ! *i*) = *Some* (*ACs* ! *i*))
unfolding *FunA RuleC* **using** *join-rule-fun RuleC*(*3*) *join-sym* **by** *metis*
have *those* (*map2* (⊔) (*map* τ (*var-rule* α)) *BCs*) ≠ *None* **proof**−
{**fix** i **assume** i:i < *length* (*zip* (*map* τ (*var-rule* α)) *BCs*)
from i **obtain** x **where** x:*var-rule* α ! i = x x ∈ *vars-term* (*to-pterm* (*lhs* α))
by (*metis* (*no-types*, *lifting*) *comp-apply length-map length-zip min-less-iff-conj nth-mem set-remdups set-rev set-vars-term-list vars-to-pterm*)
have τ (*var-rule* α ! i) ⊲ A **proof**−
from *RuleC*(*2*) *no-var-lhs* **obtain** f' *ts* **where** *lhs* α = *Fun* f' *ts* **by** *fastforce*
with x **show** *?thesis*
using *subst-image-subterm*[*of* x] *match-lhs-A* **unfolding** *FunA*
by (*smt* (*verit*) *match-matches to-pterm.simps*(*2*))
qed
moreover have τ (*var-rule* α ! i) ∈ *wf-pterm R*
using i *match-well-def*[*OF subterm*(*2*) *match-lhs-A*] **by** (*simp add*: *vars-to-pterm*)
moreover have σ (*var-rule* α ! i) ∈ *wf-pterm R*
using i *match-well-def*[*OF subterm*(*3*) *match-lhs-B*] **by** (*simp add*: *vars-to-pterm*)
moreover have τ (*var-rule* α ! i) ⊔ σ (*var-rule* α ! i) ≠ *None*
using *join-pterm-subst-Some x match-lhs-B match-lhs-A match-matches subterm.prems*(*4*) x **by** *blast*

moreover have τ (*var-rule* α ! *i*) $\sqcup$ (*Cs*!*i*) $\neq$ *None*
using *args-AC* **by** (*metis join-sym RuleC*(*3*) *i length-map length-zip min-less-iff-conj option.distinct*(*1*))
moreover have σ (*var-rule* α ! *i*) $\sqcup$ (*Cs*!*i*) $\neq$ *None*
using *args-BC* **by** (*metis join-sym RuleC*(*3*) *i length-map length-zip min-less-iff-conj option.distinct*(*1*))
ultimately have *IH*:*join-opt* (τ (*var-rule* α ! *i*)) (σ (*var-rule* α ! *i*) $\sqcup$ (*Cs*!*i*)) $\neq$ *None*
using *RuleC*(*3*,*4*) *subterm*(*1*) *i* **by** *simp*
from *IH* **have** (τ (*var-rule* α ! *i*)) $\sqcup$ (*BCs*!*i*) $\neq$ *None*
using *i args-BC l-BC-C join-sym* **by** (*metis* (*no-types*, *opaque-lifting*) *join-opt.simps*(*1*) *length-zip min-less-iff-conj*)
then have (*map2* ($\sqcup$) (*map* τ (*var-rule* α)) *BCs*)!*i* $\neq$ *None*
unfolding *nth-map*[*OF i*] **using** *i* **by** *auto*
}
then show *?thesis* **by** (*simp add*: *list-all-length those-not-none-xs*)
qed
with *match-lhs-A* **show** *?thesis*
unfolding *joinBC BC FunA* **unfolding** *fg join-opt.simps join.simps*(*7*) **by** *force*
qed (*simp add*:*FunB*)
next
case (*RuleB* α *Bs*)
from *joinAB* **have** $*$:*Prule* α *Bs* $\sqcup$ *Pfun f As* = *Some AB* **unfolding** *FunA RuleB* **using** *join-sym* **by** *metis*
obtain σ *ABs* **where** *match-lhs-A*:*match A* (*to-pterm* (*lhs* α)) = *Some* σ
and *AB*:*AB* = *Prule* α *ABs* **and** *l-A-AB*:*length ABs* = *length Bs*
and *args-AB*:($\forall$ *i*<*length Bs. Bs* ! *i* $\sqcup$ σ (*var-rule* α ! *i*) = *Some* (*ABs* ! *i*))
unfolding *FunA RuleB* **using** *join-rule-fun*[*OF* $*$ *subterm*(*3*)[*unfolded FunA RuleB*]] *RuleB*(*3*) **by** *fastforce*
from *subterm*(*4*,*7*) **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*FunC Cs g*)
from *joinBC* **have** $*$:*Prule* α *Bs* $\sqcup$ *Pfun g Cs* = *Some BC* **unfolding** *FunC RuleB* **by** *metis*
from *subterm*(*3*) **obtain** τ *BCs* **where** *match-lhs-C*:*match C* (*to-pterm* (*lhs* α)) = *Some* τ
and *BC*:*BC* = *Prule* α *BCs* **and** *l-BC-B*:*length BCs* = *length Bs*
and *args-BC*:($\forall$ *i*<*length Bs. Bs* ! *i* $\sqcup$ τ (*var-rule* α ! *i*) = *Some* (*BCs* ! *i*))
unfolding *FunC RuleB* **using** *join-rule-fun*[*OF joinBC*[*unfolded FunC RuleB*]] *RuleB*(*3*) **by** *fastforce*
have *those* (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *BCs*) $\neq$ *None* **proof**−
{fix *i* **assume** *i*:*i* < *length* (*zip* (*map* τ (*var-rule* α)) *BCs*)
from *i* **obtain** *x* **where** *x*:*var-rule* α ! *i* = *x x* $\in$ *vars-term* (*to-pterm* (*lhs* α))
by (*metis* (*no-types*, *lifting*) *comp-apply length-map length-zip min-less-iff-conj nth-mem set-remdups set-rev set-vars-term-list vars-to-pterm*)
have σ (*var-rule* α ! *i*) $\lhd$ *A* **proof**−
from *RuleB*(*2*) *no-var-lhs* **obtain** f' *ts* **where** *lhs* α = *Fun* f' *ts* **by**

fastforce
with x **show** *?thesis*
using *subst-image-subterm*[*of* x] *match-lhs-A* **unfolding** *FunA*
by (*smt* (*verit*) *match-matches to-pterm.simps*(*2*))
qed
moreover have σ (*var-rule* α ! i) $\in$ *wf-pterm R*
using i *match-well-def*[*OF subterm*(*2*) *match-lhs-A*] **by** (*simp add*: *vars-to-pterm*)
moreover have τ (*var-rule* α ! i) $\in$ *wf-pterm R*
using i *match-well-def*[*OF subterm*(*4*) *match-lhs-C*] **by** (*simp add*: *vars-to-pterm*)
moreover have σ (*var-rule* α ! i) $\sqcup$ τ (*var-rule* α ! i) $\neq$ *None*
using *join-pterm-subst-Some x match-lhs-C match-lhs-A match-matches subterm.prems*(*6*) x **by** *blast*
moreover have (*Bs*!i) $\sqcup$ τ (*var-rule* α ! i) $\neq$ *None*
using *args-BC i* **by** (*metis RuleB*(*3*) *i length-map length-zip min-less-iff-conj option.distinct*(*1*))
moreover have σ (*var-rule* α ! i) $\sqcup$ (*Bs*!i) $\neq$ *None*
using *args-AB* **by** (*metis join-sym RuleB*(*3*) *i length-map length-zip min-less-iff-conj option.distinct*(*1*))
moreover have σ (*var-rule* α ! i) $\sqcup$ τ (*var-rule* α ! i) $\neq$ *None*
using *join-pterm-subst-Some x match-lhs-C match-lhs-A match-matches subterm.prems*(*6*) x **by** *blast*
ultimately have *IH*:*join-opt* (σ (*var-rule* α ! i)) ((*Bs*!i) $\sqcup$ τ (*var-rule* α ! i)) $\neq$ *None*
using *RuleB*(*3*,*4*) *subterm*(*1*) i **by** *simp*
then have (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *BCs*)!i $\neq$ *None*
using i *args-BC l-BC-B* **unfolding** *nth-map*[*OF i*] **using** i **by** *auto*
}
then show *?thesis* **by** (*simp add*: *list-all-length those-not-none-xs*)
qed
with *match-lhs-A* **show** *?thesis*
unfolding *joinBC BC FunA* **unfolding** *join-opt.simps join.simps*(*7*) **by** *force*
next
case (*RuleC* β *Cs*)
obtain *BCs* **where** α:α = β **and** *l-B-C*:*length Bs* = *length Cs*
and *BC*:*BC* = *Prule* α *BCs* **and** *l-BC-B*:*length BCs* = *length Bs* **and** *args-BC*:($\forall$ *i*<*length Bs. Bs* ! i $\sqcup$ *Cs* ! i = *Some* (*BCs* ! i))
using *join-rule-rule joinBC subterm*(*3*,*4*) **unfolding** *RuleB*(*1*) *RuleC*(*1*) **by** *blast*
from *joinAC match-lhs-A* **have** *args-AC*:$\forall$ *i*<*length Cs. Cs* ! i $\sqcup$ σ (*var-rule* α ! i) $\neq$ *None*
using *join-rule-fun* **by** (*metis* (*no-types*, *lifting*) *FunA*(*1*) *Residual-Join-Deletion.join-sym RuleC*(*1*) α *option.distinct*(*1*) *option.inject subterm.prems*(*3*))
have *those* (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *BCs*) $\neq$ *None* **proof**−
{fix i **assume** i:i < *length* (*zip* (*map* σ (*var-rule* α)) *BCs*)
from i **obtain** x **where** x:*var-rule* α ! i = x x $\in$ *vars-term* (*to-pterm* (*lhs* α))

by (*metis* (*no-types*, *lifting*) *comp-apply length-map length-zip min-less-iff-conj nth-mem set-remdups set-rev set-vars-term-list vars-to-pterm*)
have σ (*var-rule* α ! i) $\lhd$ A **proof**−
from *RuleB*(*2*) *no-var-lhs* **obtain** f' *ts* **where** *lhs* α = *Fun* f' *ts* **by** *fastforce*
with x **show** *?thesis*
using *subst-image-subterm*[*of* x] *match-lhs-A* **unfolding** *FunA*
by (*smt* (*verit*) *match-matches to-pterm.simps*(*2*))
qed
moreover have σ (*var-rule* α ! i) $\in$ *wf-pterm R*
using i *match-well-def*[*OF subterm*(*2*) *match-lhs-A*] **by** (*simp add*: *vars-to-pterm*)
moreover have σ (*var-rule* α ! i) $\sqcup$ (*Bs*!i) $\neq$ *None*
using *args-AB* **by** (*metis join-sym RuleB*(*3*) *i length-map length-zip min-less-iff-conj option.distinct*(*1*))
moreover have (*Bs*!i) $\sqcup$ (*Cs*!i) $\neq$ *None*
using *args-BC i* **by** (*simp add*: *l-BC-B*)
moreover have σ (*var-rule* α ! i) $\sqcup$ (*Cs*!i) $\neq$ *None*
using *args-AC* **by** (*metis join-sym RuleC*(*3*) α *i length-map length-zip min-less-iff-conj*)
ultimately have *IH*:*join-opt* (σ (*var-rule* α ! i)) ((*Bs*!i) $\sqcup$ (*Cs*!i)) $\neq$ *None*
using *RuleB*(*3*,*4*) *RuleC*(*3*,*4*) *subterm*(*1*) *i l-B-C* **by** *auto*
then have (*map2* ($\sqcup$) (*map* σ (*var-rule* α)) *BCs*)!i $\neq$ *None*
using *i args-BC l-BC-B* **unfolding** *nth-map*[*OF i*] **using** *i* **by** *auto*
}
then show *?thesis* **by** (*simp add*: *list-all-length those-not-none-xs*)
qed
with *match-lhs-A* **show** *?thesis*
unfolding *joinBC BC FunA* **unfolding** *join-opt.simps join.simps*(*7*) **by** *force*
qed (*simp add*: *FunA*)
qed (*simp add*: *FunA*)
next
case (*RuleA* α *As*)
from *subterm*(*3*,*5*) **show** *?thesis* **proof**(*cases B rule*:*wf-pterm.cases*[*case-names VarB FunB RuleB*])
case (*VarB x*)
from *subterm*(*2*−) **show** *?thesis*
by (*metis join-sym VarB joinBC join-opt.simps*(*1*) *join-with-source source.simps*(*1*) *source-join to-pterm.simps*(*1*))
next
case (*FunB Bs f*)
from *subterm*(*2*) **obtain** σ *ABs* **where** *match-lhs-B*:*match B* (*to-pterm* (*lhs* α)) = *Some* σ
and *AB*:*AB* = *Prule* α *ABs* **and** *l-A-AB*:*length ABs* = *length As*
and *args-AB*:($\forall$ i<*length As*. *As* ! i $\sqcup$ σ (*var-rule* α ! i) = *Some* (*ABs* ! i))
unfolding *RuleA FunB* **using** *join-rule-fun*[*OF joinAB*[*unfolded RuleA FunB*]] *RuleA*(*1*,*3*) **by** *fastforce*

from *subterm*(*4*,*6*) **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])

case (*FunC Cs g*)

from *subterm*(*2*) **obtain** τ *ACs* **where** *match-lhs-C*:*match C* (*to-pterm* (*lhs* α)) = *Some* τ

and *AC*:*AC* = *Prule* α *ACs* **and** *l-A-AC*:*length ACs* = *length As*

and *args-AC*:($\forall$ *i*<*length As*. *As* ! *i* $\sqcup$ τ (*var-rule* α ! *i*) = *Some* (*ACs* ! *i*))

unfolding *RuleA FunC* **using** *join-rule-fun*[*OF joinAC*[*unfolded RuleA FunC*]] *RuleA*(*1*,*3*) **by** *fastforce*

from *joinBC* **obtain** ϱ **where** $\forall$ *x*$\in$*vars-term* (*to-pterm* (*lhs* α)). σ *x* $\sqcup$ τ *x* = *Some* (ϱ *x*) **and** *BC* = *to-pterm* (*lhs* α) $\cdot$ ϱ

using *join-pterm-subst-Some*[*of to-pterm* (*lhs* α)] *match-lhs-C match-lhs-B* **by** (*smt* (*verit*) *match-matches*)

then obtain *BCs* **where** *args-BC*:($\forall$ *i*<*length As*. σ (*var-rule* α ! *i*) $\sqcup$ τ (*var-rule* α ! *i*) = *Some* (*BCs* ! *i*))

and *BC*:*BC* = (*to-pterm* (*lhs* α)) $\cdot$ $\langle BCs\rangle_\alpha$ **and** *l-A-BC*:*length As* = *length BCs*

using *subst-imp-mk-subst*[*of BC to-pterm* (*lhs* α)] *RuleA*(*3*)

by (*smt* (*verit*, *del-insts*) *comp-apply nth-mem set-remdups set-rev set-vars-term-list vars-to-pterm*)

from *RuleA*(*2*) *no-var-lhs* **obtain** f' *ts* **where** *lhs*:*lhs* α = *Fun* f' *ts* **by** *fastforce*

{**fix** *i* **assume** *i*:*i* < *length As*

from *i* **obtain** *x* **where** *x*:*var-rule* α ! *i* = *x* *x* $\in$ *vars-term* (*to-pterm* (*lhs* α))

by (*metis RuleA*(*3*) *comp-apply nth-mem set-remdups set-rev set-vars-term-list vars-to-pterm*)

have σ (*var-rule* α ! *i*) $\in$ *wf-pterm R*

using *RuleA*(*3*) *i match-well-def*[*OF subterm*(*3*) *match-lhs-B*] **by** (*simp add*: *vars-to-pterm*)

moreover have τ (*var-rule* α ! *i*) $\in$ *wf-pterm R*

using *RuleA*(*3*) *i match-well-def*[*OF subterm*(*4*) *match-lhs-C*] **by** (*simp add*: *vars-to-pterm*)

moreover have *As*!*i* $\sqcup$ σ (*var-rule* α ! *i*) $\neq$ *None*

using *args-AB i* **by** *auto*

moreover have *As*!*i* $\sqcup$ τ (*var-rule* α ! *i*) $\neq$ *None*

using *args-AC i* **by** *auto*

moreover have σ (*var-rule* α ! *i*) $\sqcup$ τ (*var-rule* α ! *i*) $\neq$ *None*

using *args-BC i* **by** *auto*

ultimately have *IH*:*join-opt* (*As*!*i*) (σ (*var-rule* α ! *i*) $\sqcup$ τ (*var-rule* α ! *i*)) $\neq$ *None*

using *RuleA*(*3*,*4*) *subterm*(*1*) *i* **by** (*metis RuleA*(*1*) *nth-mem supt.arg*)

then have *As*!*i* $\sqcup$ *BCs*!*i* $\neq$ *None*

using *i args-BC* **by** *auto*

}

with *subterm*(*2*) **show** *?thesis*

unfolding *joinBC BC RuleA*(*1*) **unfolding** *join-opt.simps* **using** *join-rule-lhs l-A-BC* **by** *auto*

next

case (*RuleC* β *Cs*)
from *joinBC subterm*(*4*) **obtain** τ *BCs* **where** *match-lhs-B2*:*match B* (*to-pterm* (*lhs* β)) = *Some* τ
and *BC*:*BC* = *Prule* β *BCs* **and** *l-BC-C*:*length BCs* = *length Cs*
and *args-BC*:($\forall$ *i*<*length Cs*. *Cs* ! *i* $\sqcup$ τ (*var-rule* β ! *i*) = *Some* (*BCs* ! *i*))
unfolding *FunB RuleC* **using** *join-rule-fun RuleC*(*1*,*2*,*3*) *join-sym* **by** *metis*
from *joinAC* **have** α:$\alpha = \beta$ **and** *l-A-C*:*length As* = *length Cs*
unfolding *RuleA RuleC* **by**(*metis join.simps*(*3*) *option.distinct*(*1*))+
have *those* (*map2* ($\sqcup$) *As BCs*) $\neq$ *None* **proof**−
{**fix** *i* **assume** *i*:*i* < *length* (*zip As BCs*)
moreover have τ (*var-rule* β ! *i*) $\in$ *wf-pterm R*
using *i match-well-def*[*OF subterm*(*3*) *match-lhs-B2*] **by** (*simp add*: *RuleA*(*3*) α *vars-to-pterm*)
moreover have *As*!*i* $\sqcup$ τ (*var-rule* β ! *i*) $\neq$ *None*
using *join-pterm-subst-Some match-lhs-B subterm*(*5*) α *args-AB i match-lhs-B2* **by** *auto*
moreover have (*As*!*i*) $\sqcup$ (*Cs*!*i*) $\neq$ *None*
using *joinAC RuleA*(*1*) *RuleC*(*1*) *i join-rule-rule subterm.prems*(*1*) *subterm.prems*(*3*) **by** *fastforce*
moreover have τ (*var-rule* β ! *i*) $\sqcup$ (*Cs*!*i*) $\neq$ *None*
using *args-BC i* **by** (*simp add*: *join-sym l-A-C*)
ultimately have *IH*:*join-opt* (*As*!*i*) (τ (*var-rule* β ! *i*) $\sqcup$ (*Cs*!*i*)) $\neq$ *None*

using *RuleA*(*1*,*3*,*4*) *RuleC*(*3*,*4*) *subterm*(*1*) *i* α **by** *simp*
from *IH* **have** (*As*!*i*) $\sqcup$ (*BCs*!*i*) $\neq$ *None*
using *i args-BC* **by** (*simp add*: *join-sym l-BC-C*)
then have (*map2* ($\sqcup$) *As BCs*)!*i* $\neq$ *None*
unfolding *nth-map*[*OF i*] **using** *i* **by** *auto*
}
then show *?thesis* **by** (*simp add*: *list-all-length those-not-none-xs*)
qed
then show *?thesis*
unfolding *joinBC BC RuleA* α **unfolding** *join-opt.simps join.simps*(*7*) **by** *force*
qed (*simp add*:*FunB*)
next
case (*RuleB* β *Bs*)
from *joinAB* **have** $\alpha\beta$:$\alpha = \beta$ **and** *l-A-B*:*length As* = *length Bs*
unfolding *RuleA RuleB* **by**(*metis join.simps*(*3*) *option.distinct*(*1*))+
obtain *ABs* **where** *AB*:*AB* = *Prule* α *ABs* **and** *l-AB-B*:*length ABs* = *length Bs*
and *args-AB*:($\forall$ *i*<*length ABs*. *As*!*i* $\sqcup$ *Bs*!*i* = *Some* (*ABs* ! *i*))
using *join-rule-rule*[*OF joinAB*[*unfolded RuleA RuleB*]] *subterm*(*2*,*3*) *RuleA*(*1*) *RuleB*(*1*) **by** *metis*
from *subterm*(*4*,*6*) **show** *?thesis* **proof**(*cases C rule*:*wf-pterm.cases*[*case-names VarC FunC RuleC*])
case (*VarC x*)
from *joinBC RuleB*(*2*) *no-var-lhs* **show** *?thesis* **unfolding** *VarC RuleB*

by (*metis Residual-Join-Deletion.join-sym RuleB*(*1*) *VarC join-opt.simps*(*1*) *join-with-source source.simps*(*1*) *source-join subterm.prems*(*2*) *subterm.prems*(*3*) *subterm.prems*(*4*) *to-pterm.simps*(*1*))
next
case (*FunC Cs f*)
from *subterm*(*3*) **obtain** σ *BCs* **where** *match-lhs-C*:*match C* (*to-pterm* (*lhs* β)) = *Some* σ
and *BC*:*BC* = *Prule* β *BCs* **and** *l-BC-C*:*length BCs* = *length Bs*
and *args-BC*:($\forall$ *i*<*length Bs. Bs* ! *i* $\sqcup$ σ (*var-rule* β ! *i*) = *Some* (*BCs* ! *i*))
unfolding *RuleB* **using** *join-rule-fun*[*OF joinBC*[*unfolded RuleB FunC*]] *RuleB*(*1*,*2*,*3*) **by** (*metis FunC*(*1*))
have *those* (*map2* ($\sqcup$) *As BCs*) $\neq$ *None* **proof**−
{**fix** *i* **assume** *i*:*i* < *length* (*zip As BCs*)
have σ (*var-rule* β ! *i*) $\in$ *wf-pterm R*
using *i match-well-def*[*OF subterm*(*4*) *match-lhs-C*] **by** (*simp add*: *RuleA*(*3*) $\alpha\beta$ *vars-to-pterm*)
moreover have *As*!*i* $\sqcup$ σ (*var-rule* β ! *i*) $\neq$ *None*
using *match-lhs-C joinAC* $\alpha\beta$ *args-AB i* **unfolding** *RuleA*(*1*) *FunC*
by (*metis* (*no-types*, *lifting*) *RuleA*(*1*) *join-rule-fun length-zip min-less-iff-conj option.distinct*(*1*) *option.sel subterm.prems*(*1*))
ultimately have *IH*:*join-opt* (*As*!*i*) ((*Bs*!*i*) $\sqcup$ (σ (*var-rule* β ! *i*))) $\neq$ *None*
using *RuleA*(*1*,*3*,*4*) *subterm*(*1*) *i args-AB args-BC*
by (*metis* (*no-types*, *lifting*) *RuleB*(*4*) *l-AB-B l-A-B length-zip min-less-iff-conj nth-mem option.distinct*(*1*) *supt.arg*)
from *IH* **have** (*As*!*i*) $\sqcup$ (*BCs*!*i*) $\neq$ *None*
using *i args-BC* **by** (*simp add*: *join-sym l-BC-C*)
then have (*map2* ($\sqcup$) *As BCs*)!*i* $\neq$ *None*
unfolding *nth-map*[*OF i*] **using** *i* **by** *auto*
}
then show *?thesis* **by** (*simp add*: *list-all-length those-not-none-xs*)
qed
then show *?thesis*
unfolding *joinBC BC RuleA* $\alpha\beta$ **unfolding** *join-opt.simps join.simps*(*7*) **by** *force*
next
case (*RuleC* γ *Cs*)
from *joinBC* **have** $\beta\gamma$:$\beta = \gamma$ **and** *l-B-C*:*length Bs* = *length Cs*
using *RuleB RuleC join-rule-rule* **by** *blast*+
obtain *BCs* **where** *BC*:*BC* = *Prule* β *BCs* **and** *l-BC-B*:*length BCs* = *length Bs*
and *args-BC*:($\forall$ *i*<*length BCs. Bs*!*i* $\sqcup$ *Cs*!*i* = *Some* (*BCs* ! *i*))
using *join-rule-rule*[*OF joinBC*[*unfolded RuleB RuleC*]] *subterm*(*3*,*4*) *RuleB*(*1*) *RuleC*(*1*) **by** *metis*
obtain *ACs* **where** *AC*:*AC* = *Prule* α *ACs* **and** *l-AC-C*:*length ACs* = *length Cs*
and *args-AC*:($\forall$ *i*<*length ACs. As*!*i* $\sqcup$ *Cs*!*i* = *Some* (*ACs* ! *i*))
using *join-rule-rule*[*OF joinAC*[*unfolded RuleA RuleC*]] *subterm*(*2*,*4*) *RuleA*(*1*) *RuleC*(*1*) **by** *metis*

have *those* (*map2* (⊔) *As BCs*) ≠ *None* **proof**−
{**fix** *i* **assume** *i*:*i* < *length* (*zip As BCs*)
from *RuleA*(*1*,*4*) *RuleB*(*1*,*4*) *RuleC*(*1*,*4*) *i* **have** *join-opt* (*As*!*i*) ((*Bs*!*i*) ⊔ (*Cs*!*i*)) ≠ *None*
using *subterm*(*1*) *l-A-B l-B-C l-AC-C l-AB-B args-AB args-AC args-BC*
by (*smt* (*verit*) *length-zip min-less-iff-conj nth-mem option.distinct*(*1*) *supt.arg*)
then have (*map2* (⊔) *As BCs*)!*i* ≠ *None*
using *i args-BC* **by** *simp*
}
then show *?thesis*
by (*simp add*: *list-all-length those-not-none-xs*)
qed
then show *?thesis*
unfolding *joinBC BC* **unfolding** *RuleA αβ join-opt.simps* **by** (*simp add*: *option.case-eq-if*)
qed
qed
qed
qed

lemma *join-list-defined*:
assumes ∀ *a1 a2*. *a1* ∈ *set As* ∧ *a2* ∈ *set As* ⟶ *a1* ⊔ *a2* ≠ *None*
and ∀ *a* ∈ *set As*. *a* ∈ *wf-pterm R* **and** *As* ≠ []
shows ∃ *D*. *join-list As* = *Some D* ∧ *D* ∈ *wf-pterm R*
using *assms* **proof**(*induct length As arbitrary*:*As rule*:*full-nat-induct*)
case *1*
then show *?case* **proof**(*cases As rule*:*list.exhaust*[*case-names empty As*])
case (*As A1 As′*)
with *1* **show** *?thesis* **proof**(*cases As′ rule*:*list.exhaust*[*case-names empty As′*])
case (*As′ A2 As″*)
have *A1-wf*:*A1* ∈ *wf-pterm R* **and** *A2-wf*:*A2* ∈ *wf-pterm R*
using *1*(*3*) **unfolding** *As As′* **by** *auto*
from *As′ 1*(*2*) **obtain** *A12* **where** *A12*:*A1* ⊔ *A2* = *Some A12*
unfolding *As* **by** *fastforce*
with *A1-wf A2-wf* **have** *A12-wf*:*A12* ∈ *wf-pterm R*
by (*simp add*: *join-wf-pterm*)
show *?thesis* **proof**(*cases As″* = [])
case *True*
show *?thesis*
unfolding *As As′ True join-list.simps* **using** *A12 A12-wf* **by** *simp*
next
case *False*
from *1* **obtain** *D′* **where** *D′*:*join-list As″* = *Some D′ D′* ∈ *wf-pterm R*
unfolding *As As′* **by** (*metis False Suc-le-length-iff impossible-Cons list.set-intros*(*2*) *nat-le-linear*)
from *1*(*2*) **have** ∀ *a1 a2*. *a1* ∈ *set* (*A2* # *As″*) ∧ *a2* ∈ *set* (*A2* # *As″*) ⟶ *a1* ⊔ *a2* ≠ *None*
unfolding *As As′* **by** *force*

moreover have *Suc* (*length* (*A2* # *As''*)) $\leq$ *length As*
unfolding *As As'* **by** *simp*
moreover have ($\forall$ *a*$\in$*set* (*A2* # *As''*). *a* $\in$ *wf-pterm R*)
using *1*(*3*) **unfolding** *As As'* **by** *simp*
moreover have *A2* # *As''* $\neq$ [] **by** *simp*
ultimately obtain *D* **where** *join-list* (*A2* # *As''*) = *Some D* **and** *D-wf*:*D* $\in$ *wf-pterm R*
using *1*(*1*) **by** *blast*
then have *D*:*A2* $\sqcup$ *D'* = *Some D*
using *D' False* **using** *join-list.elims* **by** *force*
moreover have *A1* $\sqcup$ *D'* $\neq$ *None* **proof**−
from *1*(*2*) **have** $\forall$ *a1 a2*. *a1* $\in$ *set* (*A1* # *As''*) $\wedge$ *a2* $\in$ *set* (*A1* # *As''*) $\longrightarrow$ *a1* $\sqcup$ *a2* $\neq$ *None*
unfolding *As As'* **by** *force*
moreover have *Suc* (*length* (*A1* # *As''*)) $\leq$ *length As*
unfolding *As As'* **by** *simp*
moreover have ($\forall$ *a*$\in$*set* (*A1* # *As''*). *a* $\in$ *wf-pterm R*)
using *1*(*3*) **unfolding** *As As'* **by** *simp*
moreover have *A1* # *As''* $\neq$ [] **by** *simp*
ultimately have *join-list* (*A1* # *As''*) $\neq$ *None*
using *1*(*1*) **by** (*metis option.simps*(*3*))
with *D'* **show** *?thesis*
by (*metis False join-list.simps*(*3*) *join-opt.simps*(*1*) *list.exhaust*)
qed
moreover have *A1* $\sqcup$ *A2* $\neq$ *None*
using *1*(*2*) **unfolding** *As As'* **by** *simp*
ultimately have *join-opt A1* (*A2* $\sqcup$ *D'*) $\neq$ *None*
using *join-triple-defined D' A1-wf A2-wf* **unfolding** *join-list.simps* **by** *blast*
moreover have *join-list As* = *join-opt A1* (*A2* $\sqcup$ *D'*)
unfolding *As As'* **using** *False* **by** (*metis D'*(*1*) *join-list.simps*(*3*) *join-opt.simps*(*1*) *neq-Nil-conv*)
ultimately show *?thesis*
unfolding *D join-opt.simps* **using** *D-wf A1-wf join-wf-pterm* **by** *fastforce*
qed
qed *simp*
qed *simp*
qed

lemma *join-list-wf-pterm*:
assumes $\forall$ *a* $\in$ *set As*. *a* $\in$ *wf-pterm R*
and *join-list As* = *Some B*
shows *B* $\in$ *wf-pterm R*
using *assms* **proof**(*induct As arbitrary*:*B*)
case (*Cons A As*)
then show *?case* **proof**(*cases As* = [])
case *True*
from *Cons*(*2*,*3*) **show** *?thesis* **unfolding** *join-list.simps True* **by** *simp*
next

case *False*
with *Cons*(*3*) **obtain** B' **where** B':*join-list As* = *Some* B'
by (*smt* (*verit*, *ccfv-threshold*) *join-list.elims join-opt.elims list.inject*)
with *Cons* **have** $B' \in$ *wf-pterm R*
by *simp*
then show *?thesis* **using** B' *Cons*
by (*metis False join-list.simps*(*3*) *join-opt.simps*(*1*) *join-wf-pterm list.set-intros*(*1*) *neq-Nil-conv*)
qed
qed *simp*

lemma *source-join-list*:
assumes *join-list As* = *Some B* **and** $\forall$ *a* $\in$ *set As*. *a* $\in$ *wf-pterm R*
shows $\bigwedge$*A*. *A* $\in$ *set As* $\Longrightarrow$ *source A* = *source B*
proof−
fix *Ai* **assume** *Ai* $\in$ *set As*
then show *co-initial Ai B* **using** *assms* **proof**(*induct As arbitrary*: *B*)
case *Nil*
then show *?case* **by** (*simp add*: *source-join*)
next
case (*Cons A As*)
show *?case* **proof**(*cases As* = [])
case *True*
from *Cons* **show** *?thesis* **unfolding** *True*
by (*simp add*: *source-join*)
next
case *False*
have *wf*:*A* $\in$ *wf-pterm R* $\forall$ *a*$\in$*set As*. *a* $\in$ *wf-pterm R*
using *Cons*(*4*) **by** *simp-all*
from *Cons*(*2*,*3*)**obtain** B' **where** B':*join-list As* = *Some* B' *join-list* (*A*#*As*) = *join-opt A* (*Some* B')
by (*metis False join-list.simps*(*3*) *join-opt.simps*(*2*) *list.exhaust option.exhaust*)
show *?thesis* **proof**(*cases Ai* = *A*)
case *True*
show *?thesis* **unfolding** *True*
using B' *Cons*(*3*) *False source-join wf* **by** (*metis join-list-wf-pterm join-opt.simps*(*1*))
next
case *False*
then have *Ai* $\in$ *set As*
using *Cons*(*2*) **by** *simp*
with *Cons*(*1*) B'(*1*) *wf*(*2*) **have** *co-initial Ai* B'
by *simp*
moreover from B'(*1*) *wf* **have** $B' \in$ *wf-pterm R*
using *join-list-wf-pterm* **by** *blast*
ultimately show *?thesis*
by (*metis* B'(*2*) *Cons.prems*(*2*) *Residual-Join-Deletion.join-sym join-opt.simps*(*1*) *local.wf*(*1*) *source-join*)
qed

qed
qed
qed

end

3.3 Deletion

fun *deletion* :: $('f,\ 'v)$ *pterm* $\Rightarrow$ $('f,'v)$ *pterm* $\Rightarrow$ $('f,'v)$ *pterm option* (**infixr** $-_p$ *70*)
where
Var x $-_p$ *Var y* =
(*if* $x = y$ *then Some* (*Var x*) *else None*)
| *Pfun f As* $-_p$ *Pfun g Bs* =
(*if* ($f = g \wedge$ *length As* = *length Bs*) *then*
(*case those* (*map2* ($-_p$) *As Bs*) *of*
Some xs $\Rightarrow$ *Some* (*Pfun f xs*)
| *None* $\Rightarrow$ *None*)
else None)
| *Prule* α *As* $-_p$ *Prule* β *Bs* =
(*if* $\alpha = \beta$ *then*
(*case those* (*map2* ($-_p$) *As Bs*) *of*
Some xs $\Rightarrow$ *Some* ((*to-pterm* (*lhs* α)) $\cdot$ $\langle xs\rangle_\alpha$)
| *None* $\Rightarrow$ *None*)
else None)
| *Prule* α *As* $-_p$ *B* =
(*case match B* (*to-pterm* (*lhs* α)) *of*
None $\Rightarrow$ *None*
| *Some* σ $\Rightarrow$
(*case those* (*map2* ($-_p$) *As* (*map* σ (*var-rule* α))) *of*
Some xs $\Rightarrow$ *Some* (*Prule* α *xs*)
| *None* $\Rightarrow$ *None*))
| *A* $-_p$ *B* = *None*

lemma *del-empty*:
assumes $A \in$ *wf-pterm R*
shows *A* $-_p$ (*to-pterm* (*source A*)) = *Some A*
using *assms* **proof** (*induction A*)
case (*2 As f*)
then have *those* (*map2 deletion As* (*map* (*to-pterm* $\circ$ *source*) *As*)) = *Some As*
by (*simp add*:*those-some*)
then show *?case* **by** *simp*
next
case (*3* α *As*)
then have σ: *match* (*to-pterm* (*lhs* α $\cdot$ $\langle$*map source As*$\rangle_\alpha$)) (*to-pterm* (*lhs* α)) = *Some* ($\langle$*map* (*to-pterm* $\circ$ *source*) *As*$\rangle_\alpha$)
by (*metis* (*no-types*, *lifting*) *fun-mk-subst lhs-subst-trivial map-map to-pterm.simps*(*1*) *to-pterm-subst*)
from *3* **have** *those* (*map2 deletion As* (*map* (*to-pterm* $\circ$ *source*) *As*)) = *Some*

```
As
    by (simp add:those-some)
  then have args:those (map2 deletion As (map (⟨map (to-pterm ∘ source) As⟩_α)
(var-rule α))) = Some As
    by (metis 3.hyps(2) apply-lhs-subst-var-rule length-map)
  show ?case proof(cases source (Prule α As))
    case (Var x)
    then show ?thesis
      using σ residual.simps(4)[of α As x] args by auto
  next
    case (Fun f ts)
    then show ?thesis
      using σ residual.simps(5)[of α As f] args by auto
  qed
qed simp

context no-var-lhs
begin
lemma deletion-source:
assumes A ∈ wf-pterm R B ∈ wf-pterm R
    and A −_p B = Some C
  shows source C = source A
  using assms proof(induct A arbitrary:B C)
  case (1 x)
  then show ?case proof (cases B)
    case (1 y)
    then show ?thesis
      by (metis 1.prems(2) deletion.simps(1) option.distinct(1) option.inject)
  next
    case (3 α As)
    with 1 no-var-lhs show ?thesis
      by simp
  qed simp
next
  case (2 As f)
  then show ?case proof(cases B)
    case (Pfun g Bs)
    from 2(3) have f:f = g
      unfolding Pfun by (metis deletion.simps(2) not-None-eq)
    from 2(3) have l:length As = length Bs
      unfolding Pfun by (metis deletion.simps(2) not-None-eq)
    from 2(3) obtain Cs where cs:those (map2 (−_p) As Bs) = Some Cs
      unfolding Pfun f using l by fastforce
    with 2(3) have c:C = Pfun g Cs
      unfolding Pfun by (simp add: f l)
    from cs l have l-cs:length Cs = length As
      using length-those by force
    {fix i assume i:i < length As
      with 2(2) have Bs!i ∈ wf-pterm R
```

by (*metis Pfun fun-well-arg l nth-mem*)
moreover from *2*(*3*) *i cs* **have** $As!i -_p Bs!i = Some\ (Cs!i)$
using *l those-some2* **by** *fastforce*
ultimately have *source* $(Cs!i)$ = *source* $(As!i)$
using *2*(*1*) **using** *i nth-mem* **by** *blast*
}
then show *?thesis* **unfolding** *c f*
using *l-cs* **by** (*simp add*: *map-nth-eq-conv*)
qed *simp-all*
next
case (*3* α *As*)
from *3*(*1*) *no-var-lhs* **obtain** *f ts* **where** *f*:*lhs* α = *Fun f ts*
by *blast*
then show *?case* **proof**(*cases B*)
case (*Var x*)
have *match* (*Var x*) (*to-pterm* (*lhs* α)) = *None*
unfolding *f* **by** (*smt* (*verit*, *ccfv-SIG*) *Term.term.simps*(*4*) *match-matches not-Some-eq source.simps*(*1*) *source-to-pterm subst-apply-eq-Var*)
with *3*(*5*) **show** *?thesis*
unfolding *Var* **using** *f deletion.simps*(*4*) **by** *simp*
next
case (*Pfun g Bs*)
from *3*(*5*) **obtain** σ **where** *sigma*:*match B* (*to-pterm* (*lhs* α)) = *Some* σ
unfolding *Pfun* **using** *deletion.simps*(*5*) **by** *fastforce*
with *3*(*5*) **obtain** *Cs* **where** *cs*:*those* (*map2* $(-_p)$ *As* (*map* σ (*var-rule* α))) = *Some Cs*
unfolding *Pfun* **by** *fastforce*
with *3*(*5*) **have** *c*:*C* = *Prule* α *Cs*
using *sigma* **unfolding** *Pfun* **by** *simp*
from *cs 3*(*2*) **have** *l-cs*:*length Cs* = *length As*
using *length-those* **by** *force*
{fix *x* **assume** $x \in$ *vars-term* (*lhs* α)
then obtain *i* **where** *i*:$i <$ *length* (*var-rule* α) *var-rule* α $!i = x$
by (*metis in-set-conv-nth set-vars-term-list vars-term-list-vars-distinct*)
then have σ (*var-rule* α ! *i*) $\in$ *wf-pterm R*
using *match-well-def*[*OF 3*(*4*) *sigma*] **by** (*metis vars-to-pterm*)
moreover from *i cs* **have** $As!i -_p \sigma$ (*var-rule* α ! *i*) = *Some* $(Cs!i)$
using *those-some2 3.hyps*(*2*) **by** *fastforce*
ultimately have *source* $(Cs!i)$ = *source* $(As!i)$
using *3*(*3*) **using** *i nth-mem 3.hyps*(*2*) **by** *force*
then have *source* $((\langle As\rangle_\alpha)\ x)$ = *source* $((\langle Cs\rangle_\alpha)\ x)$ **using** *i*
by (*metis 3.hyps*(*2*) *l-cs lhs-subst-var-i*)
}
then show *?thesis* **unfolding** *c* **using** *l-cs 3*(*2*) **unfolding** *source.simps*
by (*smt* (*verit*, *best*) *apply-lhs-subst-var-rule comp-def in-set-conv-nth length-map nth-map set-remdups set-rev set-vars-term-list term-subst-eq-conv*)
next
case (*Prule* β *Bs*)
from *3*(*5*) **have** *alpha*:$\alpha = \beta$

unfolding *Prule* **by** (*metis deletion.simps*(*3*) *option.distinct*(*1*))
with *3* **have** *l*:*length As* = *length Bs*
unfolding *Prule* **using** *wf-pterm.cases* **by** *force*
from *3*(*5*) **obtain** *Cs* **where** *cs*:*those* (*map2* ($-_p$) *As Bs*) = *Some Cs*
unfolding *Prule alpha* **by** *fastforce*
with *3*(*5*) **have** *c*:*C* = *to-pterm* (*lhs* α) $\cdot$ $\langle Cs\rangle_\alpha$
unfolding *Prule alpha* **by** *simp*
from *cs l* **have** *l-cs*:*length Cs* = *length As*
using *length-those* **by** *force*
{**fix** *i* **assume** *i*:*i* < *length As*
with *3*(*4*) **have** *Bs*!*i* $\in$ *wf-pterm R*
unfolding *Prule* **by** (*metis fun-well-arg l nth-mem*)
moreover from *i cs* **have** *As*!*i* $-_p$ *Bs*!*i* = *Some* (*Cs*!*i*)
using *l those-some2* **by** *fastforce*
ultimately have *source* (*Cs*!*i*) = *source* (*As*!*i*)
using *3*(*3*) **using** *i nth-mem* **by** *blast*
}
then show *?thesis*
unfolding *c* **using** *l-cs* **unfolding** *source.simps* **using** *source-apply-subst*
by (*metis fun-mk-subst nth-map-conv source.simps*(*1*) *source-to-pterm to-pterm-wf-pterm*)
qed
qed
end

3.4 Computations With Single Redexes

When a proof term contains only a single rule symbol, we say it is a **single redex*.

definition *ll-single-redex* :: ('*f*, '*v*) *term* ⇒ *pos* ⇒ ('*f*, '*v*) *prule* ⇒ ('*f*, '*v*) *pterm*
where *ll-single-redex s p* α = (*ctxt-of-pos-term p* (*to-pterm s*))⟨*Prule* α (*map* (*to-pterm* ∘ (λ*pi*. *s*|-(*p*@*pi*))) (*var-poss-list* (*lhs* α)))⟩

The *ll* in *ll-single-redex* stands for **left−linear*, since this definition only makes sense for left-linear rules.

lemma *source-single-redex*:
assumes *p* ∈ *poss s*
shows *source* (*ll-single-redex s p* α) = (*ctxt-of-pos-term p s*)⟨(*lhs* α) $\cdot$ ⟨*map* (λ*pi*. *s*|-(*p*@*pi*)) (*var-poss-list* (*lhs* α))⟩$_\alpha$⟩
proof−
have *source* (*Prule* α (*map* (*to-pterm* ∘ (λ*pi*. *s*|-(*p*@*pi*))) (*var-poss-list* (*lhs* α)))) = (*lhs* α) $\cdot$ ⟨*map* (λ*pi*. *s*|-(*p*@*pi*)) (*var-poss-list* (*lhs* α))⟩$_\alpha$
unfolding *source.simps* **using** *map-nth-eq-conv* **by** *fastforce*
with *assms* **show** *?thesis*
unfolding *ll-single-redex-def* **by** (*metis context-source source-to-pterm to-pterm-ctxt-of-pos-apply-term*)

qed

lemma *target-single-redex*:

assumes $p \in$ *poss s*
shows *target* (*ll-single-redex s p* α) = (*ctxt-of-pos-term p s*)$\langle$(*rhs* α) $\cdot$ $\langle$*map* (λ*pi*. *s*|-(*p*@*pi*)) (*var-poss-list* (*lhs* α))$\rangle_\alpha\rangle$
proof−
have *target* (*Prule* α (*map* (*to-pterm* $\circ$ (λ*pi*. *s*|-(*p*@*pi*))) (*var-poss-list* (*lhs* α)))) = (*rhs* α) $\cdot$ $\langle$*map* (λ*pi*. *s*|-(*p*@*pi*)) (*var-poss-list* (*lhs* α))$\rangle_\alpha$
unfolding *target.simps* **by** (*metis* (*no-types*, *lifting*) *fun-mk-subst map-map target-empty-apply-subst target-to-pterm to-pterm.simps*(*1*) *to-pterm-empty to-pterm-subst*)
with *assms* **show** *?thesis*
unfolding *ll-single-redex-def* **using** *target-to-pterm-ctxt to-pterm-ctxt-at-pos* **by** *metis*
qed

lemma *single-redex-rstep*:
assumes *to-rule* $\alpha \in R$ $p \in$ *poss s*
shows (*source* (*ll-single-redex s p* α), *target* (*ll-single-redex s p* α)) $\in$ *rstep R*
using *source-single-redex target-single-redex assms* **by** *blast*

lemma *single-redex-neq*:
assumes (α, *p*) $\neq$ (β, *q*) $p \in$ *poss s* $q \in$ *poss s*
shows *ll-single-redex s p* α $\neq$ *ll-single-redex s q* β
proof−
from *assms*(*1*) **consider** $\alpha \neq \beta \wedge p = q \mid p \neq q$
by *fastforce*
then show *?thesis* **proof**(*cases*)
case *1*
then have *Prule* α (*map* (*to-pterm* $\circ$ (λ*pi*. *s* |- (*p* @ *pi*))) (*var-poss-list* (*lhs* α))) $\neq$ *Prule* β (*map* (*to-pterm* $\circ$ (λ*pi*. *s* |- (*p* @ *pi*))) (*var-poss-list* (*lhs* α)))
by *simp*
with *1* **show** *?thesis*
using *assms*(*2*,*3*) **unfolding** *ll-single-redex-def* **by** *simp*
next
case *2*
show *?thesis* **proof**(*cases* $p \in$ *poss* (*ll-single-redex s q* β))
case *True*
from *2* **consider** (*qp*) $q <_p p$ | (*par*) $q \perp p$ | (*pq*) $p <_p q$
using *pos-cases* **by** *force*
then show *?thesis* **proof**(*cases*)
case *qp*
then obtain *i r* **where** *r*:*q*@(*i*#*r*) = *p*
using *less-pos-def′* **by** (*metis neq-Nil-conv*)
with ‹$p \in$ *poss* (*ll-single-redex s q* β)› **have** *i*#*r* $\in$ *poss* (*Prule* β (*map* (*to-pterm* $\circ$ (λ*pi*. *s* |- (*q* @ *pi*))) (*var-poss-list* (*lhs* β))))
unfolding *ll-single-redex-def* **using** *assms*(*3*) **by** (*metis hole-pos-ctxt-of-pos-term hole-pos-poss-conv poss-list-sound poss-list-to-pterm*)
then have *i*:*i* < *length* (*var-poss-list* (*lhs* β)) **and** *r* $\in$ *poss* (*map* (*to-pterm* $\circ$ (λ*pi*. *s* |- (*q* @ *pi*))) (*var-poss-list* (*lhs* β))!*i*)
by *auto*
then have *r* $\in$ *poss* (*to-pterm* (*s* |- (*q* @ (*var-poss-list* (*lhs* β)!*i*))))

by *simp*
then obtain s' **where** (*Prule* β (*map* (*to-pterm* $\circ$ (λ*pi*. *s* |- (*q* @ *pi*))) (*var-poss-list* (*lhs* β))))|-(*i*#*r*) = *to-pterm* s'
by (*metis* (*no-types*, *lifting*) *comp-apply ctxt-supt-id i nth-map poss-list-sound poss-list-to-pterm subt-at.simps*(*2*) *subt-at-hole-pos to-pterm-ctxt-of-pos-apply-term*)

then have (*Prule* β (*map* (*to-pterm* $\circ$ (λ*pi*. *s* |- (*q* @ *pi*))) (*var-poss-list* (*lhs* β))))|-(*i*#*r*) $\neq$ *Prule* α (*map* (*to-pterm* $\circ$ (λ*pi*. *s* |- (*p* @ *pi*))) (*var-poss-list* (*lhs* α)))
using *to-pterm.elims* **by** *auto*
then show *?thesis* **using** *r assms*(*2*,*3*) **unfolding** *ll-single-redex-def*
by (*smt* (*verit*, *del-insts*) *hole-pos-ctxt-of-pos-term hole-pos-poss p-in-poss-to-pterm replace-at-subt-at subt-at-append*)
next
case *par*
then have *ll-single-redex s q* β |- *p* = *to-pterm s* |-*p*
using *True* **unfolding** *ll-single-redex-def*
by (*simp add*: *assms*(*2*,*3*) *p-in-poss-to-pterm parallel-replace-at-subt-at*)
then show *?thesis*
using *assms*(*2*,*3*) **unfolding** *ll-single-redex-def*
by (*metis ctxt-supt-id hole-pos-ctxt-of-pos-term is-empty-step.simps*(*3*) *p-in-poss-to-pterm subt-at-hole-pos to-pterm-ctxt-of-pos-apply-term to-pterm-empty*)
next
case *pq*
then obtain *r* **where** *r*:*q* = *p*@*r* *r* $\neq$ []
using *less-pos-def'* **by** *blast*
then have $*$:*ll-single-redex s q* β |- *p* = (*ctxt-of-pos-term r* (*to-pterm* (*s*|-*p*)))⟨*Prule* β (*map* (*to-pterm* $\circ$ (λ*pi*. *s* |- (*q* @ *pi*))) (*var-poss-list* (*lhs* β)))⟩
using *True* **unfolding** *ll-single-redex-def r* **by** (*metis* (*no-types*, *lifting*) *assms*(*2*) *ctxt-apply-subt-at ctxt-supt-id p-in-poss-to-pterm replace-at-subt-at to-pterm-ctxt-of-pos-apply-term*)
from *r*(*2*) *assms*(*3*) **obtain** *f ts i r'* **where** *f*:*s*|-*p* = *Fun f ts* **and** *r'*:*r* = *i*#*r'*
unfolding *r* **by** (*metis args-poss neq-Nil-conv poss-append-poss*)
have *ll-single-redex s q* β |- *p* $\neq$ *Prule* α (*map* (*to-pterm* $\circ$ (λ*pi*. *s* |- (*p* @ *pi*))) (*var-poss-list* (*lhs* α)))
unfolding $*$ **unfolding** *ll-single-redex-def f to-pterm.simps r' ctxt-of-pos-term.simps intp-actxt.simps* **by** *simp*
then show *?thesis*
using *assms*(*2*) **unfolding** *ll-single-redex-def*
by (*metis p-in-poss-to-pterm replace-at-subt-at*)
qed
next
case *False*
then show *?thesis* **unfolding** *ll-single-redex-def* **using** *assms*(*2*)
by (*metis hole-pos-ctxt-of-pos-term hole-pos-poss p-in-poss-to-pterm*)
qed
qed
qed

context *left-lin-wf-trs*
begin
lemma *rstep-exists-single-redex*:
  assumes $(s, t) \in$ *rstep R*
  shows $\exists$ *A p* α. *A* = (*ll-single-redex s p* α) $\wedge$ *source A* = *s* $\wedge$ *target A* = *t* $\wedge$
to-rule $\alpha \in R \wedge p \in$ *poss s*
proof−
  from *assms* **obtain** *C* σ *l r* **where** *lr*:$(l, r) \in R$ **and** *s*:$s = C\langle l \cdot \sigma\rangle$ **and** *t*: $t = C\langle r \cdot \sigma\rangle$
    by *blast*
  from *s* **obtain** *p* **where** *p*:$p \in$ *poss s* **and** *C*:*C* = *ctxt-of-pos-term p s*
    using *hole-pos-poss* **by** *fastforce*
  let *?subst*=⟨*map* (λ*pi. s* |- (*p* @ *pi*)) (*var-poss-list l*)⟩$_{(}$*l* →*r*)
  {**fix** *x* **assume** $x \in$ *vars-term l*
    then obtain *i* **where** *i*:*i* < *length* (*vars-term-list l*) *vars-term-list l* ! *i* = *x*
      by (*metis in-set-idx set-vars-term-list*)
    with *left-lin lr* **have** *var-l*:*vars-distinct l* ! *i* = *x*
      using *linear-term-var-vars-term-list left-linear-trs-def* **by** *fastforce*
    let *?p*=*var-poss-list l* ! *i*
    from *i* **have** *l* |- *?p* = *Var x* **using** *vars-term-list-var-poss-list* **by** *auto*
    moreover have $l \cdot \sigma$ = *s*|-*p* **using** *s C p replace-at-subt-at* **by** *fastforce*
    ultimately have *left*:σ *x* = (*s* |-*p*) |- *?p*
      by (*metis eval-term.simps*(*1*) *i*(*1*) *length-var-poss-list nth-mem subt-at-subst*
var-poss-imp-poss var-poss-list-sound)
    from *i* **have** *map* (λ*pi. s* |- (*p* @ *pi*)) (*var-poss-list l*) !*i* = (*s* |-*p*) |- *?p*
      by (*simp add*: *length-var-poss-list p*)
    with *left var-l* **have** σ *x* = *?subst x* **unfolding** *mk-subst-def prule.sel*
     by (*smt* (*verit*, *best*) *case-prodE comp-apply distinct-rev i*(*1*) *left-lin left-linear-trs-def*
length-map length-var-poss-list linear-term-var-vars-term-list lr mk-subst-def mk-subst-distinct
prod.sel(*1*) *remdups-id-iff-distinct rev-rev-ident*)
  }**note** *subst*=*this*
   then have (*ctxt-of-pos-term p s*)⟨*l* · ⟨*map* (λ*pi. s* |- (*p* @ *pi*)) (*var-poss-list*
l)⟩$_{(}$*l*→*r*)⟩⟩ = $C\langle l \cdot \sigma\rangle$
    using *C* **by** (*simp add*: *eval-same-vars*)
  then have *source* (*ll-single-redex s p* (*Rule l r*)) = *s*
    using *source-single-redex*[*OF p*] *s* **by** *auto*
  moreover have *target* (*ll-single-redex s p* (*l* → *r*)) = *t*
   using *subst varcond lr target-single-redex*[*OF p*] *eval-same-vars-cong* **unfolding**
t C
    by (*smt* (*verit*) *case-prodD prule.sel*(*1*) *prule.sel*(*2*) *vars-term-subset-subst-eq*)
  ultimately show *?thesis* **using** *lr p* **by** *fastforce*
qed
end

lemma *single-redex-wf-pterm*:
  assumes *to-rule* $\alpha \in R$ **and** *lin*:*linear-term* (*lhs* α)
    and *p*:$p \in$ *poss s*
  shows *ll-single-redex s p* $\alpha \in$ *wf-pterm R*

proof−
from *lin* **have** *l*:*length* (*map* (*to-pterm* ∘ (*λpi. s* |- (*p* @ *pi*))) (*var-poss-list* (*lhs* *α*))) = *length* (*var-rule* *α*)
using *length-var-poss-list linear-term-var-vars-term-list* **by** *fastforce*
have *Prule* *α* (*map* (*to-pterm* ∘ (*λpi. s* |- (*p* @ *pi*))) (*var-poss-list* (*lhs* *α*))) ∈ *wf-pterm R*
using *wf-pterm.intros*(*3*)[*OF assms*(*1*) *l*] *to-pterm-wf-pterm* **by** *force*
then show *?thesis* **unfolding** *ll-single-redex-def*
using *ctxt-wf-pterm p to-pterm-wf-pterm* **by** (*metis p-in-poss-to-pterm*)
qed

Interaction of a single redex Δ, contained in A with the proof term A.

locale *single-redex* = *left-lin-no-var-lhs* +
fixes *A* Δ *p q α*
assumes *a-well*:*A* ∈ *wf-pterm R*
and *p*:*p* ∈ *poss* (*source A*) **and** *q*:*q* ∈ *poss A*
and *pq*:*ctxt-of-pos-term p* (*source A*) = *source-ctxt* (*ctxt-of-pos-term q A*)
and *delta*:Δ = *ll-single-redex* (*source A*) *p α*
and *aq*:*A*|-*q* = *Prule α* (*map* (*λi. A*|-(*q*@[*i*])) [*0*..<*length* (*var-rule α*)])
begin

interpretation *residual-op*:*op-proof-term R residual*
using *op-proof-term.intro*[*OF left-lin-no-var-lhs-axioms*] *op-proof-term-axioms.intro*[*of R residual*] *res-empty2* **by** *force*

interpretation *deletion-op*:*op-proof-term R deletion*
using *op-proof-term.intro*[*OF left-lin-no-var-lhs-axioms*] *op-proof-term-axioms.intro*[*of R deletion*] *del-empty* **by** *force*

abbreviation *As* ≡ (*map* (*λi. A*|-(*q*@[*i*])) [*0*..<*length* (*var-rule α*)])

lemma *length-as*:*length As* = *length* (*var-rule α*)
by *simp*

lemma *as-well*:∀ *i* < *length As*. *As*!*i* ∈ *wf-pterm R*
using *subt-at-is-wf-pterm a-well aq q*
by (*metis fun-well-arg nth-mem*)

lemma *a*:*A* = (*ctxt-of-pos-term q A*)⟨*Prule α As*⟩
using *aq* **by** (*simp add*: *q replace-at-ident*)

lemma *rule-in-TRS*: *to-rule α* ∈ *R*
proof−
from *a-well a q* **have** *Prule α As* ∈ *wf-pterm R*
by (*metis subt-at-ctxt-of-pos-term subt-at-is-wf-pterm*)
then show *?thesis*
using *wf-pterm.cases* **by** *force*
qed

lemma *lin-lhs*:*linear-term* (*lhs* α)
  using *rule-in-TRS left-lin left-linear-trs-def* **by** *fastforce*

lemma *source-at-pq*:*source* (*A*|-*q*) = (*source A*)|-*p*
proof−
  from *a-well q* **have** (*ctxt-of-pos-term q A*) $\in$ *wf-pterm-ctxt R*
    by (*simp add*: *ctxt-of-pos-term-well*)
  then have *source A* = (*source-ctxt* (*ctxt-of-pos-term q A*)) $\langle$*source* (*A*|-*q*)$\rangle$
    using *source-ctxt-apply-term q* **by** (*metis ctxt-supt-id*)
  moreover from *p* **have** *source A* = (*ctxt-of-pos-term p* (*source A*)) $\langle$(*source A*)|-*p*$\rangle$
    by (*simp add*: *replace-at-ident*)
  ultimately show *?thesis*
    using *pq p q* **by** *simp*
qed

lemma *single-redex-pterm*:
  shows Δ = (*ctxt-of-pos-term p* (*to-pterm* (*source A*)))$\langle$*Prule* α (*map* (*to-pterm* $\circ$ *source*) *As*)$\rangle$
proof−
  from *lin-lhs* **have** *l2*:*length* (*var-poss-list* (*lhs* α)) = *length* (*var-rule* α)
      by (*metis length-var-poss-list linear-term-var-vars-term-list*)
  {**fix** *i* **assume** *i*:*i* < *length* (*var-poss-list* (*lhs* α))
    let *?pi*=*var-poss-list* (*lhs* α)!*i*
    from *i* **have** $*$:(*lhs* α)|-*?pi* = *Var* ((*var-rule* α)!*i*)
        using *lin-lhs* **by** (*metis linear-term-var-vars-term-list length-var-poss-list vars-term-list-var-poss-list*)
    from *source-at-pq* **have** *source A* |- (*p* @ *?pi*) = *source* (*Prule* α *As*)|-*?pi*
      by (*metis a p q subt-at-append subt-at-ctxt-of-pos-term*)
    also have ... = *Var* ((*var-rule* α)!*i*) $\cdot$ $\langle$*map source As*$\rangle_\alpha$
       unfolding *source.simps* **using** *subt-at-subst* $*$ *i nth-mem var-poss-imp-poss* **by** *fastforce*
    also have ... = *source* (*As*!*i*)
        unfolding *eval-term.simps* **using** *i lhs-subst-var-i length-as l2* **by** (*metis* (*no-types*, *lifting*) *length-map nth-map*)
    finally have *source A* |- (*p* @ *?pi*) = *source* (*As*!*i*) **.**
  }
  with *l2* **show** *?thesis*
    unfolding *delta ll-single-redex-def* **by** (*simp add*: *nth-map-conv*)
qed

lemma *delta-trs-wf-pterm*:
 shows Δ $\in$ *wf-pterm R*
proof−
  have *well2*:*Prule* α (*map* (*to-pterm* $\circ$ *source*) *As*) $\in$ *wf-pterm R* **proof**−
    from *a-well a q* **have** *Prule* α *As* $\in$ *wf-pterm R*
      by (*metis subt-at-ctxt-of-pos-term subt-at-is-wf-pterm*)
    then have *to-rule* α $\in$ *R*
      using *wf-pterm.cases* **by** *auto*

then show *?thesis*
by (*simp add*: *wf-pterm.intros*(*3*))
qed
show *?thesis* **unfolding** *single-redex-pterm*
using *p well2* **by** (*simp add*: *p-in-poss-to-pterm ctxt-wf-pterm*)
qed

lemma *source-delta*: *source* Δ = *source A*
proof−
have *src*:*source* (*Prule* α (*map* (*to-pterm* $\circ$ *source*) *As*)) = *source* (*Prule* α *As*)
unfolding *source.simps* **by** (*metis* (*no-types*, *lifting*) *comp-eq-dest-lhs list.map-comp list.map-cong0 source-to-pterm*)
moreover have *source-ctxt* (*ctxt-of-pos-term p* (*to-pterm* (*source A*))) = *source-ctxt* (*ctxt-of-pos-term q A*)
using *pq* **by** (*metis p source-to-pterm-ctxt′ to-pterm-ctxt-at-pos*)
ultimately show *?thesis* **unfolding** *single-redex-pterm*
using *p p q* **by** (*metis aq p-in-poss-to-pterm pq replace-at-ident source-at-pq source-ctxt-apply-term to-pterm-trs-ctxt*)
qed

lemma *residual*:
shows *A re* Δ = *Some* ((*ctxt-of-pos-term q A*)$\langle$(*to-pterm* (*rhs* α)) $\cdot$ $\langle As\rangle_\alpha$$\rangle$)
proof−
have *l*:*length* (*map2* (*re*) *As* (*map* (*to-pterm* $\circ$ *source*) *As*)) = *length As*
by *simp*
{fix *i* **assume** *i*:*i* < *length As*
with *as-well* **have** *As*!*i re* (*to-pterm* $\circ$ *source*) (*As*!*i*) = *Some* (*As*!*i*)
by (*metis* (*no-types*, *lifting*) *o-apply res-empty2*)
then have *map2* (*re*) *As* (*map* (*to-pterm* $\circ$ *source*) *As*) ! *i* = *Some* (*As* ! *i*)
using *i* **by** *force*
}
then have ∗:*those* (*map2* (*re*) *As* (*map* (*to-pterm* $\circ$ *source*) *As*)) = *Some As*
using *those-some*[*OF l*] **using** *l* **by** *presburger*
then have (*Prule* α *As*) *re* (*Prule* α (*map* (*to-pterm* $\circ$ *source*) *As*)) = *Some* ((*to-pterm* (*rhs* α)) $\cdot$ $\langle As\rangle_\alpha$)
using *residual.simps*(*3*)[*of* α *As* α (*map* (*to-pterm* $\circ$ *source*) *As*)] **by** *simp*
moreover from *single-redex-pterm* **have** Δ = (*to-pterm-ctxt* (*source-ctxt* (*ctxt-of-pos-term q A*)))$\langle$(*Prule* α (*map* (*to-pterm* $\circ$ *source*) *As*))$\rangle$
unfolding *delta ll-single-redex-def pq*[*symmetric*] **by** (*simp add*: *p to-pterm-ctxt-at-pos*)
ultimately show *?thesis*
using *a residual-op.apply-f-ctxt* **by** (*metis a-well ctxt-of-pos-term-well q*)
qed

lemma *residual-well*:
the (*A re* Δ) $\in$ *wf-pterm R*
using *a-well* **by** (*metis delta-trs-wf-pterm option.sel residual residual-well-defined*)

lemma *target-residual*:*target* (*the* (*A re* Δ)) = *target A*
apply(*subst* (*2*) *a*)

unfolding *residual option.sel*
apply(*subst* (*1 2*) *context-target*)
by (*metis fun-mk-subst target.simps*(*1*) *target.simps*(*3*) *target-empty-apply-subst target-to-pterm to-pterm-empty*)

lemma *deletion*:
shows $A -_p \Delta = Some\ ((ctxt\text{-}of\text{-}pos\text{-}term\ q\ A)\langle(to\text{-}pterm\ (lhs\ \alpha)) \cdot \langle As\rangle_\alpha\rangle)$
proof−
have *l*:$length\ (map2\ (-_p)\ As\ (map\ (to\text{-}pterm \circ source)\ As)) = length\ As$
by *simp*
{**fix** *i* **assume** *i*:$i < length\ As$
with *as-well* **have** $As!i -_p (to\text{-}pterm \circ source)\ (As!i) = Some\ (As!i)$
by (*metis* (*no-types*, *lifting*) *o-apply del-empty*)
then have $map2\ (-_p)\ As\ (map\ (to\text{-}pterm \circ source)\ As)\ !\ i = Some\ (As\ !\ i)$
using *i* **by** *force*
}
then have $*$:$those\ (map2\ (-_p)\ As\ (map\ (to\text{-}pterm \circ source)\ As)) = Some\ As$
using *those-some*[*OF l*] **using** *l* **by** *presburger*
then have $(Prule\ \alpha\ As) -_p (Prule\ \alpha\ (map\ (to\text{-}pterm \circ source)\ As)) = Some\ ((to\text{-}pterm\ (lhs\ \alpha)) \cdot \langle As\rangle_\alpha)$
using *deletion.simps*(*3*)[*of* α *As* α (*map* (*to-pterm* $\circ$ *source*) *As*)] **by** *simp*
moreover from *single-redex-pterm* **have** $\Delta = (to\text{-}pterm\text{-}ctxt\ (source\text{-}ctxt\ (ctxt\text{-}of\text{-}pos\text{-}term\ q\ A)))\langle(Prule\ \alpha\ (map\ (to\text{-}pterm \circ source)\ As))\rangle$
unfolding *delta ll-single-redex-def pq*[*symmetric*] **by** (*simp add*: *p to-pterm-ctxt-at-pos*)
ultimately show *?thesis*
using *a deletion-op.apply-f-ctxt* **by** (*metis a-well ctxt-of-pos-term-well q*)
qed

lemma *deletion-well*:
shows $the\ (A -_p \Delta) \in wf\text{-}pterm\ R$
proof−
have $\forall\ a \in set\ As.\ a \in wf\text{-}pterm\ R$
by (*metis a a-well fun-well-arg q subt-at-ctxt-of-pos-term subt-at-is-wf-pterm*)
then have $to\text{-}pterm\ (lhs\ \alpha) \cdot \langle As\rangle_\alpha \in wf\text{-}pterm\ R$
by (*meson lhs-subst-well-def nth-mem to-pterm-wf-pterm*)
then show *?thesis* **unfolding** *deletion option.sel*
by (*simp add*: *a-well ctxt-wf-pterm q*)
qed

end

locale *single-redex′* = *left-lin-wf-trs* +
fixes $A\ \Delta\ p\ q\ \alpha\ \sigma$
assumes *a-well*:$A \in wf\text{-}pterm\ R$ **and** *rule-in-TRS*:$to\text{-}rule\ \alpha \in R$
and *p*:$p \in poss\ (source\ A)$ **and** *q*:$q \in poss\ A$
and *pq*:$ctxt\text{-}of\text{-}pos\text{-}term\ p\ (source\ A) = source\text{-}ctxt\ (ctxt\text{-}of\text{-}pos\text{-}term\ q\ A)$
and *delta*:$\Delta = ll\text{-}single\text{-}redex\ (source\ A)\ p\ \alpha$
and *aq*:$A|\text{-}q = (to\text{-}pterm\ (lhs\ \alpha)) \cdot \sigma$

begin

interpretation *residual-op*:*op-proof-term R residual* **proof**−
  have ∗:*left-lin-no-var-lhs R*
    by (*simp add*: *left-lin-axioms left-lin-no-var-lhs.intro no-var-lhs-axioms*)
  then show *op-proof-term R* (*re*)
   using *op-proof-term.intro*[*OF* ∗] *op-proof-term-axioms.intro*[*of R residual*] *res-empty2*
by *force*
qed

lemma *a*:$A = (ctxt\text{-}of\text{-}pos\text{-}term\ q\ A)\langle(to\text{-}pterm\ (lhs\ \alpha)) \cdot \sigma\rangle$
  using *aq* **by** (*simp add*: *q replace-at-ident*)

lemma *lin-lhs*:*linear-term* (*lhs* α)
  using *rule-in-TRS left-lin left-linear-trs-def* **by** *fastforce*

lemma *is-fun-lhs*:*is-Fun* (*lhs* α)
  using *rule-in-TRS* **using** *no-var-lhs* **by** *blast*

abbreviation $As \equiv map\ \sigma\ (var\text{-}rule\ \alpha)$

lemma *lhs-subst*: $(to\text{-}pterm\ (lhs\ \alpha)) \cdot \sigma = (to\text{-}pterm\ (lhs\ \alpha)) \cdot \langle As\rangle_\alpha$
proof−
  {**fix** x **assume** $x \in vars\text{-}term\ (to\text{-}pterm\ (lhs\ \alpha))$
    then obtain i **where** $x = var\text{-}rule\ \alpha!i$ **and** $i < length\ (var\text{-}rule\ \alpha)$
     by (*metis in-set-idx set-vars-term-list vars-term-list-vars-distinct vars-to-pterm*)
    then have $\sigma\ x = (\langle As\rangle_\alpha)\ x$
      by (*metis* (*mono-tags*, *lifting*) *apply-lhs-subst-var-rule length-map nth-map*)
  }
  then show *?thesis*
    using *term-subst-eq-conv* **by** *blast*
qed

lemma *rhs-subst*: $(to\text{-}pterm\ (rhs\ \alpha)) \cdot \sigma = (to\text{-}pterm\ (rhs\ \alpha)) \cdot \langle As\rangle_\alpha$
proof−
  {**fix** x **assume** $x \in vars\text{-}term\ (to\text{-}pterm\ (rhs\ \alpha))$
    then have $x \in vars\text{-}term\ (to\text{-}pterm\ (lhs\ \alpha))$
      using *no-var-lhs varcond rule-in-TRS set-vars-term-list subsetD vars-to-pterm*
by (*metis case-prodD*)
    then obtain i **where** $x = var\text{-}rule\ \alpha!i$ **and** $i < length\ (var\text{-}rule\ \alpha)$
     by (*metis in-set-idx set-vars-term-list vars-term-list-vars-distinct vars-to-pterm*)
    then have $\sigma\ x = (\langle As\rangle_\alpha)\ x$
      by (*metis* (*mono-tags*, *lifting*) *apply-lhs-subst-var-rule length-map nth-map*)
  }
  then show *?thesis*
    using *term-subst-eq-conv* **by** *blast*
qed

lemma *as-well*:$\forall\, i < length\ As.\ As!i \in wf\text{-}pterm\ R$

using *a-well aq* **by** (*metis length-map lhs-subst lhs-subst-args-wf-pterm nth-mem q subt-at-is-wf-pterm*)

lemma *source-at-pq*:*source* (A|-q) = (*source* A)|-p
proof−
from *a-well q* **have** (*ctxt-of-pos-term q A*) ∈ *wf-pterm-ctxt R*
by (*simp add*: *ctxt-of-pos-term-well*)
then have *source A* = (*source-ctxt* (*ctxt-of-pos-term q A*)) ⟨*source* (A|-q)⟩
using *source-ctxt-apply-term q* **by** (*metis ctxt-supt-id*)
moreover from *p* **have** *source A* = (*ctxt-of-pos-term p* (*source A*)) ⟨(*source A*)|-p⟩
by (*simp add*: *replace-at-ident*)
ultimately show *?thesis*
using *pq p q* **by** *simp*
qed

lemma *single-redex-pterm*:
shows Δ = (*ctxt-of-pos-term p* (*to-pterm* (*source A*)))⟨*Prule* α (*map* (*to-pterm* ∘ *source*) *As*)⟩
proof−
from *lin-lhs* **have** *l2*:*length* (*var-poss-list* (*lhs* α)) = *length* (*var-rule* α)
by (*metis length-var-poss-list linear-term-var-vars-term-list*)
{**fix** *i* **assume** *i*:i < *length* (*var-poss-list* (*lhs* α))
let *?pi*=*var-poss-list* (*lhs* α)!i
from *i* **have** ∗:(*lhs* α)|-*?pi* = *Var* ((*var-rule* α)!i)
using *lin-lhs* **by** (*metis linear-term-var-vars-term-list length-var-poss-list vars-term-list-var-poss-list*)
from *source-at-pq* **have** *source A* |- (p @ *?pi*) = *source* ((*to-pterm* (*lhs* α)) · ⟨*As*⟩$_\alpha$)|-*?pi*
using *lhs-subst* **by** (*metis a p q subt-at-append subt-at-ctxt-of-pos-term*)
also have ... = *Var* ((*var-rule* α)!i) · ⟨*map source As*⟩$_\alpha$
using *subt-at-subst* ∗
by (*metis* (*no-types*, *lifting*) *fun-mk-subst i nth-mem source.simps*(*1*) *source-apply-subst source-to-pterm to-pterm-wf-pterm var-poss-imp-poss var-poss-list-sound*)
also have ... = *source* (*As*!i)
unfolding *eval-term.simps* **using** *i lhs-subst-var-i l2* **by** (*metis* (*no-types*, *lifting*) *length-map nth-map*)
finally have *source A* |- (p @ *?pi*) = *source* (*As*!i) **.**
}
with *l2* **show** *?thesis*
unfolding *delta ll-single-redex-def* **by** (*simp add*: *map-eq-conv′*)
qed

lemma *residual*:
shows *A re* Δ = *Some* ((*ctxt-of-pos-term q A*)⟨(*to-pterm* (*rhs* α)) · σ⟩)
proof−
have *l*:*length* (*map2* (*re*) *As* (*map* (*to-pterm* ∘ *source*) *As*)) = *length As*
by *simp*
{**fix** *i* **assume** *i*:i < *length As*

with *as-well* **have** *As!i re* (*to-pterm* ∘ *source*) (*As!i*) = *Some* (*As!i*)
  by (*metis comp-apply res-empty2*)
then have *map2* (*re*) *As* (*map* (*to-pterm* ∘ *source*) *As*) ! *i* = *Some* (*As* ! *i*)
  using *i* **by** *force*
}
then have ∗:*those* (*map2* (*re*) *As* (*map* (*to-pterm* ∘ *source*) *As*)) = *Some As*
  using *those-some*[*OF l*] **using** *l* **by** *presburger*
from *is-fun-lhs* **obtain** *f As′* **where** *f*:(*to-pterm* (*lhs* α) · $\langle As\rangle_\alpha$) = *Pfun f As′*
  by *fastforce*
then have *match*:*match* (*Pfun f As′*) (*to-pterm* (*lhs* α)) = *Some* ($\langle As\rangle_\alpha$)
  by (*metis lhs-subst-trivial*)
have *map*:*map* ($\langle As\rangle_\alpha$) (*var-rule* α) = *As*
  using *apply-lhs-subst-var-rule length-map* **by** *blast*
have ((*to-pterm* (*lhs* α)) · σ) *re* (*Prule* α (*map* (*to-pterm* ∘ *source*) *As*)) = *Some* ((*to-pterm* (*rhs* α)) · σ)
  unfolding *rhs-subst lhs-subst* **using** *residual.simps*(*7*)[*of f As′* α (*map* (*to-pterm* ∘ *source*) *As*)]
  unfolding *match f* **using** ∗ *map* **by** (*metis option.simps*(*5*))
moreover from *single-redex-pterm* **have** Δ = (*to-pterm-ctxt* (*source-ctxt* (*ctxt-of-pos-term q A*)))⟨(*Prule* α (*map* (*to-pterm* ∘ *source*) *As*))⟩
  unfolding *delta ll-single-redex-def pq*[*symmetric*] **by** (*simp add*: *p to-pterm-ctxt-at-pos*)
ultimately show *?thesis*
  using *a residual-op.apply-f-ctxt* **by** (*metis a-well ctxt-of-pos-term-well q*)
qed

end

end

4 Orthogonal Proof Terms

theory *Orthogonal-PT*
imports
  Residual-Join-Deletion
begin

inductive *orthogonal*::(′*f*, ′*v*) *pterm* ⇒ (′*f*, ′*v*) *pterm* ⇒ *bool* (**infixl** $\perp_p$ *50*)
  where
  Var x $\perp_p$ *Var x*
| *length As* = *length Bs* ⟹ ∀ *i* < *length As*. *As!i* $\perp_p$ *Bs!i* ⟹ *Fun f As* $\perp_p$ *Fun f Bs*
| *length As* = *length Bs* ⟹ ∀ (*a*,*b*) ∈ *set*(*zip As Bs*). *a* $\perp_p$ *b* ⟹ (*Prule* α *As*) $\perp_p$ (*to-pterm* (*lhs* α)) · $\langle Bs\rangle_\alpha$
| *length As* = *length Bs* ⟹ ∀ (*a*,*b*) ∈ *set*(*zip As Bs*). *a* $\perp_p$ *b* ⟹ (*to-pterm* (*lhs* α)) · $\langle As\rangle_\alpha$ $\perp_p$ (*Prule* α *Bs*)
lemmas *orthogonal.intros*[*intro*]

lemma *orth-symp*: *symp* ($\perp_p$)
proof

{**fix** *A B*::('*f*, '*v*) *pterm* **assume** $A \perp_p B$
then show $B \perp_p A$ **proof**(*induct*)
case (*3 As Bs* α)
then show *?case* **using** *orthogonal.intros*(*4*)[**where** α=α **and** *Bs*=*As* **and** *As*=*Bs*]
using *zip-symm* **by** *fastforce*
next
case (*4 As Bs* α)
then show *?case* **using** *orthogonal.intros*(*3*)[**where** α=α **and** *As*=*Bs* **and** *Bs*=*As*]
using *zip-symm* **by** *fastforce*
qed (*simp-all add*:*orthogonal.intros*)
}
qed

lemma *equal-imp-orthogonal*:
shows $A \perp_p A$
by(*induct A*) (*simp-all add*: *orthogonal.intros*)

lemma *source-orthogonal*:
assumes *source A* = *t*
shows $A \perp_p$ *to-pterm t*
using *assms* **proof**(*induct A arbitrary*:*t*)
case (*Prule* α *As*)
then have *t*:*to-pterm t* = (*to-pterm* (*lhs* α)) · $\langle$*map* (*to-pterm* ∘ *source*) *As*$\rangle_\alpha$
by (*metis fun-mk-subst list.map-comp source.simps*(*3*) *to-pterm.simps*(*1*) *to-pterm-subst*)
from *Prule*(*1*) **have** ∀ (*a*,*b*) ∈ *set* (*zip As* (*map* (*to-pterm* ∘ *source*) *As*)). $a \perp_p b$
by (*metis* (*mono-tags*, *lifting*) *case-prod-beta'* *comp-def in-set-zip nth-map zip-fst*)
with *t* **show** *?case*
using *orthogonal.intros*(*3*) **by** (*metis length-map*)
qed (*simp-all add*:*orthogonal.intros*)

lemma *orth-imp-co-initial*:
assumes $A \perp_p B$
shows *co-initial A B*
using *assms* **proof**(*induct rule*: *orthogonal.induct*)
case (*2 As Bs f*)
show *?case* **proof**(*cases f*)
case (*Inr g*)
with *2* **show** *?thesis* **unfolding** *Inr*
by (*simp add*: *nth-map-conv*)
next
case (*Inl* α)
with *2* **show** *?thesis* **unfolding** *Inl*
by (*metis nth-map-conv source.simps*(*3*))
qed
next
case (*3 As Bs* α)

then have *l*:*length* (*zip As Bs*) = *length As*
by *simp*
with *3* **have** *IH*:$\forall$ *i* < *length As*. *source* (*As*!*i*) = *source* (*Bs*!*i*)
by (*metis* (*mono-tags*, *lifting*) *case-prod-conv nth-mem nth-zip*)
have *src*:*source* ((*to-pterm* (*lhs* α)) · $\langle$*Bs*$\rangle_\alpha$) = (*lhs* α) · $\langle$*map source Bs*$\rangle_\alpha$
by (*simp add*: *source-apply-subst*)
from *3*(*1*) *l IH* **show** *?case* **unfolding** *src source.simps*
by (*metis nth-map-conv*)
next
case (*4 As Bs* α)
then have *l*:*length* (*zip As Bs*) = *length As*
by *simp*
with *4* **have** *IH*:$\forall$ *i* < *length As*. *source* (*As*!*i*) = *source* (*Bs*!*i*)
by (*metis* (*mono-tags*, *lifting*) *case-prod-conv nth-mem nth-zip*)
have *src*:*source* ((*to-pterm* (*lhs* α)) · $\langle$*As*$\rangle_\alpha$) = (*lhs* α) · $\langle$*map source As*$\rangle_\alpha$
by (*simp add*: *source-apply-subst*)
from *4*(*1*) *l IH* **show** *?case* **unfolding** *src source.simps*
by (*metis nth-map-conv*)
qed *simp*

If two proof terms are orthogonal then residual and join are well-defined.

lemma *orth-imp-residual-defined*:
assumes *varcond*:$\bigwedge$*l r*. (*l*, *r*) $\in$ *R* $\Longrightarrow$ *is-Fun l* $\bigwedge$*l r*. (*l*, *r*) $\in$ *S* $\Longrightarrow$ *is-Fun l*
and *A* $\perp_p$ *B*
and *A* $\in$ *wf-pterm R* **and** *B* $\in$ *wf-pterm S*
shows *A re B* $\neq$ *None*
using *assms*(*3*−) **proof**(*induct*)
case (*2 As Bs f*)
from *2*(*3*) **have** *wellAs*:$\forall$ *a* $\in$ *set As*. *a* $\in$ *wf-pterm R*
by *blast*
from *2*(*4*) **have** *wellBs*:$\forall$ *b* $\in$ *set Bs*. *b* $\in$ *wf-pterm S*
by *blast*
from *2*(*1*,*2*) *wellAs wellBs* **have** *c*:$\forall$ *i* < *length As*. ($\exists$ *C*. *As*!*i re Bs*!*i* = *Some C*)
by *auto*
from *2*(*1*) **have** *l*:*length As* = *length* (*map2* (*re*) *As Bs*)
by *simp*
from *2*(*1*) **have** $\forall$ *i* < *length As*. *As*!*i re Bs*!*i* = (*map2* (*re*) *As Bs*)!*i*
by *simp*
with *c* **obtain** *Cs* **where** $\forall$ *i* < *length As*. *As*!*i re Bs*!*i* = *Some* (*Cs*!*i*) **and** *length Cs* = *length As*
using *exists-some-list l* **by** (*metis* (*no-types*, *lifting*))
with *2* **have** ∗:*those* (*map2* (*re*) *As Bs*) = *Some Cs*
by (*simp add*: *those-some*)
show *?case* **proof**(*cases f*)
case (*Inr g*)
show *?thesis* **unfolding** *Inr residual.simps 2*(*1*) ∗ **by** *simp*
next
case (*Inl* α)

show *?thesis* **unfolding** *Inl residual.simps 2(1)* ∗ **by** *simp*
qed
next
case (*3 As Bs* α)
from *3(3) varcond* **obtain** *g ts* **where** *g*:*lhs* α = *Fun g ts*
by (*metis Inl-inject is-Fun-Fun-conv sum.simps(4) term.distinct(1) term.sel(2) wf-pterm.cases*)
then have ∗:*to-pterm* (*lhs* α) · $\langle Bs\rangle_\alpha$ = *Pfun g* (*map* (λ*t*. *t* · $\langle Bs\rangle_\alpha$) (*map to-pterm ts*))
by *simp*
from *3(3)* **have** *l1*:*length As* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
from *3(3)* **have** *wellAs*:∀ *a* ∈ *set As*. *a* ∈ *wf-pterm R*
by *blast*
from *3(1,4) l1* **have** *wellBs*:∀ *b* ∈ *set Bs*. *b* ∈ *wf-pterm S*
by (*simp add*: *lhs-subst-args-wf-pterm*)
from *3(1)* **have** *l2*:*length* (*zip As Bs*) = *length As*
by *simp*
with *3(1,2) wellAs wellBs* **have** ∀ *i* < *length As*. *As* ! *i re Bs* ! *i* ≠ *None*
by (*metis* (*mono-tags*, *lifting*) *case-prod-conv nth-mem nth-zip*)
then have *c*:∀ *i* < *length As*. (∃ *C*. *As*!*i re Bs*!*i* = *Some C*)
by *blast*
from *3(1)* **have** ∀ *i* < *length As*. *As*!*i re Bs*!*i* = (*map2* (*re*) *As Bs*)!*i*
by *simp*
with *c* **obtain** *Cs* **where** ∀ *i* < *length As*. *As*!*i re Bs*!*i* = *Some* (*Cs*!*i*) **and** *length Cs* = *length As*
using *exists-some-list l2* **by** (*metis* (*no-types*, *lifting*) *length-map*)
with *3* **have** *cs*:*those* (*map2* (*re*) *As Bs*) = *Some Cs*
by (*simp add*: *those-some*)
have *bs*:*match* (*to-pterm* (*lhs* α) · $\langle Bs\rangle_\alpha$) (*to-pterm* (*lhs* α)) = *Some* ($\langle Bs\rangle_\alpha$)
using *lhs-subst-trivial* **by** *blast*
then have (*map* ($\langle Bs\rangle_\alpha$) (*var-rule* α)) = *Bs*
using *3(1) l1 apply-lhs-subst-var-rule* **by** *force*
then show *?case* **using** *residual.simps(5)* **using** *bs cs g* **unfolding** ∗
by *simp*
next
case (*4 As Bs* α)
from *4(4) varcond* **obtain** *g ts* **where** *g*:*lhs* α = *Fun g ts*
by (*metis Inl-inject is-Fun-Fun-conv sum.simps(4) term.distinct(1) term.sel(2) wf-pterm.cases*)
then have ∗:*to-pterm* (*lhs* α) · $\langle As\rangle_\alpha$ = *Pfun g* (*map* (λ*t*. *t* · $\langle As\rangle_\alpha$) (*map to-pterm ts*))
by *simp*
from *4(4)* **have** *l1*:*length Bs* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
from *4(1,3) l1* **have** *wellAs*:∀ *a* ∈ *set As*. *a* ∈ *wf-pterm R*
by (*simp add*: *lhs-subst-args-wf-pterm*)
from *4(4)* **have** *wellBs*:∀ *b* ∈ *set Bs*. *b* ∈ *wf-pterm S*
by *blast*

from *4(1)* **have** *l2*:*length* (*zip As Bs*) = *length As*
by *simp*
with *4(1,2) wellAs wellBs* **have** $\forall$ *i* < *length As*. *As* ! *i re Bs* ! *i* $\neq$ *None*
by (*metis* (*mono-tags*, *lifting*) *case-prod-conv nth-mem nth-zip*)
then have *c*:$\forall$ *i* < *length As*. ($\exists$ *C*. *As*!*i re Bs*!*i* = *Some C*)
by *blast*
from *4(1)* **have** $\forall$ *i* < *length As*. *As*!*i re Bs*!*i* = (*map2* (*re*) *As Bs*)!*i*
by *simp*
with *c* **obtain** *Cs* **where** $\forall$ *i* < *length As*. *As*!*i re Bs*!*i* = *Some* (*Cs*!*i*) **and** *length Cs* = *length As*
using *exists-some-list l2* **by** (*metis* (*no-types*, *lifting*) *length-map*)
with *4* **have** *cs*:*those* (*map2* (*re*) *As Bs*) = *Some Cs*
by (*simp add*: *those-some*)
have *bs*:*match* (*to-pterm* (*lhs* α) $\cdot$ $\langle As\rangle_\alpha$) (*to-pterm* (*lhs* α)) = *Some* ($\langle As\rangle_\alpha$)
using *lhs-subst-trivial* **by** *blast*
have (*map* ($\langle As\rangle_\alpha$) (*var-rule* α)) = *As*
using *4(1) l1 apply-lhs-subst-var-rule* **by** *force*
then show *?case* **using** *residual.simps*(*7*) **using** *bs cs g* **unfolding** $*$
by *simp*
qed *simp*

lemma *orth-imp-join-defined*:
assumes *varcond*:$\bigwedge$*l r*. (*l*, *r*) $\in$ *R* $\Longrightarrow$ *is-Fun l*
and *A* $\perp_p$ *B*
and *A* $\in$ *wf-pterm R* **and** *B* $\in$ *wf-pterm R*
shows *A* $\sqcup$ *B* $\neq$ *None*
using *assms*(*2*$-$) **proof**(*induct*)
case (*2 As Bs f*)
from *2(3)* **have** *wellAs*:$\forall$ *a* $\in$ *set As*. *a* $\in$ *wf-pterm R*
by *blast*
from *2(4)* **have** *wellBs*:$\forall$ *b* $\in$ *set Bs*. *b* $\in$ *wf-pterm R*
by *blast*
from *2(1,2) wellAs wellBs* **have** *c*:$\forall$ *i* < *length As*. ($\exists$ *C*. *As*!*i* $\sqcup$ *Bs*!*i* = *Some C*)
by *auto*
from *2(1)* **have** *l*:*length As* = *length* (*map2* ($\sqcup$) *As Bs*)
by *simp*
from *2(1)* **have** $\forall$ *i* < *length As*. *As*!*i* $\sqcup$ *Bs*!*i* = (*map2* ($\sqcup$) *As Bs*)!*i*
by *simp*
with *c* **obtain** *Cs* **where** $\forall$ *i* < *length As*. *As*!*i* $\sqcup$ *Bs*!*i* = *Some* (*Cs*!*i*) **and** *length Cs* = *length As*
using *exists-some-list l* **by** (*metis* (*no-types*, *lifting*))
with *2* **have** $*$:*those* (*map2* ($\sqcup$) *As Bs*) = *Some Cs*
by (*simp add*: *those-some*)
show *?case* **proof**(*cases f*)
case (*Inr g*)
show *?thesis* **unfolding** *Inr join.simps 2(1)* $*$ **by** *simp*
next

case (*Inl* α)
show *?thesis* **unfolding** *Inl join.simps 2*(*1*) $*$ **by** *simp*
qed
next
case (*3 As Bs* α)
from *3*(*3*) *varcond* **obtain** *g ts* **where** *g*:*lhs* α = *Fun g ts*
by (*metis Inl-inject is-Fun-Fun-conv sum.simps*(*4*) *term.distinct*(*1*) *term.sel*(*2*) *wf-pterm.cases*)
then have $*$:*to-pterm* (*lhs* α) $\cdot$ $\langle Bs\rangle_\alpha$ = *Pfun g* (*map* ($\lambda t.$ *t* $\cdot$ $\langle Bs\rangle_\alpha$) (*map to-pterm ts*))
by *simp*
from *3*(*3*) **have** *l1*:*length As* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
from *3*(*3*) **have** *wellAs*:$\forall$ *a* $\in$ *set As. a* $\in$ *wf-pterm R*
by *blast*
from *3*(*1*,*4*) *l1* **have** *wellBs*:$\forall$ *b* $\in$ *set Bs. b* $\in$ *wf-pterm R*
by (*simp add*: *lhs-subst-args-wf-pterm*)
from *3*(*1*) **have** *l2*:*length* (*zip As Bs*) = *length As*
by *simp*
with *3*(*1*,*2*) *wellAs wellBs* **have** $\forall$ *i* < *length As. As* ! *i* $\sqcup$ *Bs* ! *i* $\neq$ *None*
by (*metis* (*mono-tags*, *lifting*) *case-prod-conv nth-mem nth-zip*)
then have *c*:$\forall$ *i* < *length As.* ($\exists$ *C. As*!*i* $\sqcup$ *Bs*!*i* = *Some C*)
by *blast*
from *3*(*1*) **have** $\forall$ *i* < *length As. As*!*i* $\sqcup$ *Bs*!*i* = (*map2* ($\sqcup$) *As Bs*)!*i*
by *simp*
with *c* **obtain** *Cs* **where** $\forall$ *i* < *length As. As*!*i* $\sqcup$ *Bs*!*i* = *Some* (*Cs*!*i*) **and** *length Cs* = *length As*
using *exists-some-list l2* **by** (*metis* (*no-types*, *lifting*) *length-map*)
with *3* **have** *cs*:*those* (*map2* ($\sqcup$) *As Bs*) = *Some Cs*
by (*simp add*: *those-some*)
have *bs*:*match* (*to-pterm* (*lhs* α) $\cdot$ $\langle Bs\rangle_\alpha$) (*to-pterm* (*lhs* α)) = *Some* ($\langle Bs\rangle_\alpha$)
using *lhs-subst-trivial* **by** *blast*
then have (*map* ($\langle Bs\rangle_\alpha$) (*var-rule* α)) = *Bs*
using *3*(*1*) *l1 apply-lhs-subst-var-rule* **by** *force*
then show *?case* **using** *residual.simps*(*5*) **using** *bs cs g* **unfolding** $*$
by *simp*
next
case (*4 As Bs* α)
from *4*(*4*) *varcond* **obtain** *g ts* **where** *g*:*lhs* α = *Fun g ts*
by (*metis Inl-inject is-Fun-Fun-conv sum.simps*(*4*) *term.distinct*(*1*) *term.sel*(*2*) *wf-pterm.cases*)
then have $*$:*to-pterm* (*lhs* α) $\cdot$ $\langle As\rangle_\alpha$ = *Pfun g* (*map* ($\lambda t.$ *t* $\cdot$ $\langle As\rangle_\alpha$) (*map to-pterm ts*))
by *simp*
from *4*(*4*) **have** *l1*:*length Bs* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
from *4*(*1*,*3*) *l1* **have** *wellAs*:$\forall$ *a* $\in$ *set As. a* $\in$ *wf-pterm R*
by (*simp add*: *lhs-subst-args-wf-pterm*)
from *4*(*4*) **have** *wellBs*:$\forall$ *b* $\in$ *set Bs. b* $\in$ *wf-pterm R*

by *blast*
from *4*(*1*) **have** *l2*:*length* (*zip As Bs*) = *length As*
by *simp*
with *4*(*1*,*2*) *wellAs wellBs* **have** $\forall\, i < length\ As.\ As\ !\ i \sqcup Bs\ !\ i \neq None$
by (*metis* (*mono-tags*, *lifting*) *case-prod-conv nth-mem nth-zip*)
then have *c*:$\forall\, i < length\ As.\ (\exists\, C.\ As!i \sqcup Bs!i = Some\ C)$
by *blast*
from *4*(*1*) **have** $\forall\, i < length\ As.\ As!i \sqcup Bs!i = (map2\ (\sqcup)\ As\ Bs)!i$
by *simp*
with *c* **obtain** *Cs* **where** $\forall\, i < length\ As.\ As!i \sqcup Bs!i = Some\ (Cs!i)$ **and** *length Cs* = *length As*
using *exists-some-list l2* **by** (*metis* (*no-types*, *lifting*) *length-map*)
with *4* **have** *cs*:*those* (*map2* ($\sqcup$) *As Bs*) = *Some Cs*
by (*simp add*: *those-some*)
have *bs*:*match* (*to-pterm* (*lhs* α) $\cdot$ $\langle As\rangle_\alpha$) (*to-pterm* (*lhs* α)) = *Some* ($\langle As\rangle_\alpha$)
using *lhs-subst-trivial* **by** *blast*
have (*map* ($\langle As\rangle_\alpha$) (*var-rule* α)) = *As*
using *4*(*1*) *l1 apply-lhs-subst-var-rule* **by** *force*
then show *?case* **using** *residual.simps*(*7*) **using** *bs cs g* **unfolding** $*$
by *simp*
qed *simp*

context *no-var-lhs*
begin
lemma *orth-imp-residual-defined*:
assumes $A \perp_p B$ **and** $A \in$ *wf-pterm R* **and** $B \in$ *wf-pterm R*
shows *A re B* $\neq$ *None*
using *orth-imp-residual-defined assms no-var-lhs* **by** *fastforce*

lemma *orth-imp-join-defined*:
assumes $A \perp_p B$ **and** $A \in$ *wf-pterm R* **and** $B \in$ *wf-pterm R*
shows $A \sqcup B \neq$ *None*
using *orth-imp-join-defined assms no-var-lhs* **by** *fastforce*

lemma *orthogonal-ctxt*:
assumes $C\langle A\rangle \perp_p C\langle B\rangle$ $C \in$ *wf-pterm-ctxt R*
shows $A \perp_p B$
using *assms* **proof**(*induct C*)
case (*Cfun f ss1 C ss2*)
from *Cfun*(*2*) **have** $\forall\, i{<}length\ (ss1\ @\ C\langle A\rangle\ \#\ ss2).\ (ss1\ @\ C\langle A\rangle\ \#\ ss2)\ !\ i \perp_p (ss1\ @\ C\langle B\rangle\ \#\ ss2)\ !\ i$
unfolding *intp-actxt.simps* **using** *orthogonal.simps* **by** (*smt* (*verit*) *is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *term.distinct*(*1*) *term.sel*(*4*))
then have $C\langle A\rangle \perp_p C\langle B\rangle$
by (*metis append.right-neutral length-append length-greater-0-conv length-nth-simps*(*1*) *list.discI nat-add-left-cancel-less nth-append-length*)
with *Cfun*(*1*,*3*) **show** *?case* **by** *auto*
next
case (*Crule* α *ss1 C ss2*)

from *Crule(3)* **obtain** *f ts* **where** *lhs* α = *Fun f ts*
using *no-var-lhs* **by** (*smt* (*verit*, *del-insts*) *Inl-inject Inr-not-Inl case-prodD actxt.distinct*(*1*) *actxt.inject term.collapse*(*2*) *wf-pterm-ctxt.simps*)
with *Crule(2)* **have** $\forall$ *i*<*length* (*ss1* @ $C\langle A\rangle$ # *ss2*). (*ss1* @ $C\langle A\rangle$ # *ss2*) ! *i* $\perp_p$ (*ss1* @ $C\langle B\rangle$ # *ss2*) ! *i*
unfolding *intp-actxt.simps* **using** *orthogonal.simps*
by (*smt* (*verit*, *ccfv-threshold*) *Inl-inject Inr-not-Inl eval-term.simps*(*2*) *term.distinct*(*1*) *term.inject*(*2*) *to-pterm.simps*(*2*))
then have $C\langle A\rangle \perp_p C\langle B\rangle$
by (*metis intp-actxt.simps*(*2*) *hole-pos.simps*(*2*) *hole-pos-poss nth-append-length poss-Cons-poss term.sel*(*4*))
with *Crule*(*1*,*3*) **show** *?case* **by** *auto*
qed *simp*

end

context *left-lin-no-var-lhs*
begin

lemma *orthogonal-subst*:
assumes $A \cdot \sigma \perp_p B \cdot \sigma$ *source A* = *source B*
and $A \in$ *wf-pterm R* $B \in$ *wf-pterm R*
shows $A \perp_p B$
using *assms*(*3*,*4*,*1*,*2*) **proof**(*induct A arbitrary*:*B rule*:*subterm-induct*)
case (*subterm A*)
show *?case* **proof**(*cases A*)
case (*Var x*)
with *subterm no-var-lhs* **have** *B* = *Var x*
by (*metis Inl-inject Inr-not-Inl case-prodD co-initial-Var is-VarI term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.simps*)
then show *?thesis*
unfolding *Var* **by** (*simp add*: *orthogonal.intros*(*1*))
next
case (*Pfun f As*)
with *subterm*(*5*) **show** *?thesis* **proof**(*cases B*)
case (*Pfun g Bs*)
from *subterm*(*5*) **have** *f*:*f* = *g*
unfolding ‹*A* = *Pfun f As*› *Pfun* **by** *simp*
from *subterm*(*5*) **have** *l*:*length As* = *length Bs*
unfolding ‹*A* = *Pfun f As*› *Pfun* **using** *map-eq-imp-length-eq* **by** *auto*
{**fix** *i* **assume** *i*:*i* < *length As*
with *subterm*(*4*) **have** *As*!*i* $\cdot\ \sigma \perp_p$ *Bs*!*i* $\cdot\ \sigma$
unfolding ‹*A* = *Pfun f As*› *Pfun eval-term.simps f*
by (*smt* (*verit*) *is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *length-map nth-map orthogonal.simps term.distinct*(*1*) *term.sel*(*4*))
moreover from *i subterm*(*5*) **have** *source* (*As*!*i*) = *source* (*Bs*!*i*)
unfolding ‹*A* = *Pfun f As*› *Pfun eval-term.simps f* **by** (*simp add*: *map-equality-iff*)
moreover from *i l subterm*(*2*,*3*) **have** *As*!*i* $\in$ *wf-pterm R Bs*!*i* $\in$ *wf-pterm*

```
R
        unfolding ‹A = Pfun f As› Pfun by auto
      moreover from i have As!i ◁ A
        unfolding ‹A = Pfun f As› by simp
      ultimately have As!i ⊥p Bs!i
        using subterm(1) by simp
    }
    with l show ?thesis
      unfolding f ‹A = Pfun f As› Pfun by (simp add: orthogonal.intros(2))
  next
    case (Prule β Bs)
    with subterm(3) have lin:linear-term (lhs β)
      using left-lin left-linear-trs-def wf-pterm.cases by fastforce
    from subterm(3) obtain g ts where lhs:lhs β = Fun g ts
      unfolding Prule using no-var-lhs by (metis Inl-inject case-prodD is-FunE
is-Prule.simps(1) is-Prule.simps(3) term.distinct(1) term.inject(2) wf-pterm.simps)

    with subterm(4) obtain τ1 where A · σ = to-pterm (lhs β) · τ1
        unfolding Prule Pfun eval-term.simps using orthogonal.simps by (smt
(verit, ccfv-SIG) Inl-inject Inr-not-Inl term.inject(2))
    with subterm(4,5) obtain τ where τ2:A = to-pterm (lhs β) · τ
     unfolding Prule Pfun source.simps using simple-pterm-match lin by (metis
matches-iff source.simps(2))
    let ?As=map τ (var-rule β)
    have l:length Bs = length (var-rule β)
      using Prule subterm.prems(2) wf-pterm.simps by fastforce
    from τ2 have A:A = to-pterm (lhs β) · ⟨?As⟩β
      by (metis lhs-subst-var-rule set-vars-term-list subsetI vars-to-pterm)
    {fix i assume i:i < length Bs
      have subt:?As!i ◁ A
       using i l by (metis (no-types, lifting) τ2 comp-apply lhs nth-map nth-mem
set-remdups set-rev set-vars-term-list subst-image-subterm to-pterm.simps(2) vars-to-pterm)
      have wf:?As!i ∈ wf-pterm R
         using i l by (metis A length-map lhs-subst-args-wf-pterm nth-mem sub-
term.prems(1))
      have l′:length (var-rule β) = length ?As
        by simp
      from subterm(4) A Prule lhs have orth:?As!i · σ ⊥p Bs!i · σ proof(cases)
        case (4 As′ Bs′ β′)
        then have Bs′:Bs′ = map (λs. s · σ) Bs and β′:β′ = β
          unfolding Prule by simp-all
        have As′:As′ = map (λs. s · σ) ?As proof−
         have l′′:length As′ = length ?As using 4(3) l unfolding Bs′ length-map
by simp
          {fix j assume j:j < length As′ and neq:As′!j ≠ map (λs. s · σ) ?As ! j
            let ?x=var-rule β!j
            from j have x:?x ∈ vars-term (lhs β)
         by (metis comp-apply l′ l′′ nth-mem set-remdups set-rev set-vars-term-list)
```

then obtain *p* **where** $p{:}p \in poss\ (lhs\ \beta)\ lhs\ \beta\ |\text{-}p = Var\ ?x$
by (*meson vars-term-poss-subt-at*)
from *j* **have** $1{:}(\langle As' \rangle_{\beta}{}')\ ?x = As'!j$
using $\beta'\ l''$ *lhs-subst-var-i* **by** *force*
from *j* **have** $2{:}(\langle map\ (\lambda s.\ s \cdot \sigma)\ ?As\rangle_{\beta})\ ?x = map\ (\lambda s.\ s \cdot \sigma)\ ?As\ !\ j$
using *lhs-subst-var-i* **by** (*metis* l'' *length-map*)
then have *False* **using** *4(1) 1 2 p* **unfolding** *A eval-lhs-subst*[*OF* l']
β'
by (*smt* (*verit*, *del-insts*) *x neq set-vars-term-list term-subst-eq-rev*
vars-to-pterm)
}
then show *?thesis*
using l'' **by** (*metis* (*mono-tags*, *lifting*) *map-nth-eq-conv nth-map*)
qed
have $i'{:}i < length\ (zip\ (map\ (\lambda a.\ a \cdot \sigma)\ (map\ \tau\ (var\text{-}rule\ \beta)))\ (map\ (\lambda b.$
$b \cdot \sigma)\ Bs))$
using *i l* **by** *simp*
from *4(4)* **have** $map\ (\lambda a.\ a \cdot \sigma)\ (map\ \tau\ (var\text{-}rule\ \beta))\ !\ i \perp_p map\ (\lambda b.\ b$
$\cdot\ \sigma)\ Bs\ !i$
unfolding $As'\ Bs'$ **using** i' **by** (*metis case-prodD i l length-map nth-mem*
nth-zip)
then show *?thesis*
using *i l* **by** *auto*
qed *simp-all*
have *co-init*:*source* (*?As*!*i*) = *source* (*Bs*!*i*) **proof**(*rule ccontr*)
assume *neq*:*source* (*?As*!*i*) ≠ *source* (*Bs*!*i*)
let *?x*=*var-rule* β!*i*
from *i l* **have** $x{:}?x \in vars\text{-}term\ (lhs\ \beta)$
by (*metis comp-apply nth-mem set-remdups set-rev set-vars-term-list*)
then obtain *p* **where** $p{:}p \in poss\ (lhs\ \beta)\ lhs\ \beta\ |\text{-}p = Var\ ?x$
by (*meson vars-term-poss-subt-at*)
from *i* **have** $1{:}(\langle map\ source\ ?As\rangle_{\beta})\ ?x = source\ (?As!i)$
using *l lhs-subst-var-i* **by** (*metis length-map nth-map*)
from *i* **have** $2{:}(\langle map\ source\ Bs\rangle_{\beta})\ ?x = source\ (Bs!i)$
using *l lhs-subst-var-i* **by** (*metis length-map nth-map*)
from *subterm*(*5*) **show** *False* **using** *neq 1 2 p*
unfolding *Prule A source.simps source-apply-subst*[*OF to-pterm-wf-pterm*[*of*
lhs β]] *source-to-pterm*
using *term-subst-eq-rev x* **by** *fastforce*
qed
from *subterm*(*1*) *subt wf orth co-init* **have** $?As!i \perp_p Bs!i$
using *i subterm*(*3*) **unfolding** *Prule* **by** (*meson fun-well-arg nth-mem*)
}
then show *?thesis* **unfolding** *A Prule*
by (*smt* (*verit*, *best*) *case-prodI2 in-set-idx l length-map length-zip min-less-iff-conj*
nth-zip orthogonal.intros(*4*) *prod.sel*(*1*) *snd-conv*)
qed *simp*
next
case (*Prule* α *As*)

with *subterm*(*2*) **have** *lin*:*linear-term* (*lhs* α)
using *left-lin left-linear-trs-def wf-pterm.cases* **by** *fastforce*
from *subterm*(*2*) **obtain** *f ts* **where** *lhs*:*lhs* α = *Fun f ts*
by (*metis Inr-not-Inl Prule case-prodD is-FunE no-var-lhs sum.inject*(*1*) *term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.simps*)
with *subterm*(*5*) **show** *?thesis* **proof**(*cases B*)
case (*Var x*)
then show *?thesis*
using *source-orthogonal subterm.prems*(*4*) **by** *fastforce*
next
case (*Pfun g Bs*)
with *subterm*(*4*) **obtain** $\tau 1$ **where** $B \cdot \sigma$ = *to-pterm* (*lhs* α) $\cdot$ $\tau 1$
unfolding *Prule Pfun eval-term.simps* **using** *orthogonal.simps* **by** (*smt* (*verit*, *ccfv-SIG*) *Inl-inject Inr-not-Inl term.inject*(*2*))
with *subterm*(*4*,*5*) **obtain** τ **where** $\tau 2$:B = *to-pterm* (*lhs* α) $\cdot$ τ
unfolding *Prule Pfun source.simps* **using** *simple-pterm-match lin* **by** (*metis matches-iff source.simps*(*2*))
let *?Bs*=*map* τ (*var-rule* α)
have *l*:*length As* = *length* (*var-rule* α)
using *Prule subterm.prems*(*1*) *wf-pterm.simps* **by** *fastforce*
from $\tau 2$ **have** *B*:B = *to-pterm* (*lhs* α) $\cdot$ $\langle ?Bs\rangle_\alpha$
by (*metis lhs-subst-var-rule set-vars-term-list subsetI vars-to-pterm*)
have l':*length* (*var-rule* α) = *length ?Bs*
by *simp*
{**fix** *i* **assume** *i*:i < *length As*
from *subterm*(*4*) *B Prule lhs* **have** *orth*:*As*!*i* $\cdot$ σ $\perp_p$ *?Bs*!*i* $\cdot$ σ **proof**(*cases*)
case (*3 As′ Bs′* α')
then have *As′*:*As′* = *map* ($\lambda s.\ s \cdot \sigma$) *As* **and** α':$\alpha' = \alpha$
unfolding *Prule* **by** *simp-all*
have *Bs′*:*Bs′* = *map* ($\lambda s.\ s \cdot \sigma$) *?Bs* **proof**−
have l'':*length Bs′* = *length ?Bs* **using** *3*(*3*) *l* **unfolding** *As′ length-map* **by** *simp*
{**fix** *j* **assume** *j*:j < *length Bs′* **and** *neq*:*Bs′*!*j* $\neq$ *map* ($\lambda s.\ s \cdot \sigma$) *?Bs* ! *j*
let *?x*=*var-rule* α!*j*
from *j* **have** *x*:*?x* $\in$ *vars-term* (*lhs* α)
by (*metis comp-apply* l'' *length-map nth-mem set-remdups set-rev set-vars-term-list*)
then obtain *p* **where** *p*:$p \in$ *poss* (*lhs* α) *lhs* α |-*p* = *Var ?x*
by (*meson vars-term-poss-subt-at*)
from *j* **have** *1*:($\langle Bs'\rangle_{\alpha'}$) *?x* = *Bs′*!*j*
using α' l'' *lhs-subst-var-i* **by** *force*
from *j* **have** *2*:($\langle$*map* ($\lambda s.\ s \cdot \sigma$) *?Bs*$\rangle_\alpha$) *?x* = *map* ($\lambda s.\ s \cdot \sigma$) *?Bs* ! *j*
using *lhs-subst-var-i* **by** (*metis* l'' *length-map*)
then have *False* **using** *3*(*1*) *1 2 p* **unfolding** *B eval-lhs-subst*[*OF* l'] α'
by (*smt* (*verit*, *ccfv-SIG*) *3*(*2*) *B* α' *eval-lhs-subst* l' *map-eq-conv neq set-vars-term-list term-subst-eq-rev vars-to-pterm x*)
}
then show *?thesis*

using l'' **by** (*metis* (*mono-tags*, *lifting*) *map-nth-eq-conv nth-map*)
qed
have i':i < *length* (*zip* (*map* ($\lambda b.\ b \cdot \sigma$) *As*) (*map* ($\lambda a.\ a \cdot \sigma$) (*map* τ (*var-rule* α))))
using *i l* **by** *simp*
from *3*(*4*) **have** *map* ($\lambda b.\ b \cdot \sigma$) *As* ! $i \perp_p$ *map* ($\lambda a.\ a \cdot \sigma$) (*map* τ (*var-rule* α)) ! *i*
unfolding As' Bs' **using** i' **by** (*metis case-prodD i l length-map nth-mem nth-zip*)
then show *?thesis*
using *i l* **by** *auto*
qed *simp-all*
have *co-init*:*source* (*As*!*i*) = *source* (*?Bs*!*i*) **proof**(*rule ccontr*)
assume *neq*:*source* (*As*!*i*) $\neq$ *source* (*?Bs*!*i*)
let *?x*=*var-rule* α!*i*
from *i l* **have** *x*:*?x* $\in$ *vars-term* (*lhs* α)
by (*metis comp-apply nth-mem set-remdups set-rev set-vars-term-list*)
then obtain *p* **where** *p*:*p* $\in$ *poss* (*lhs* α) *lhs* α |-*p* = *Var ?x*
by (*meson vars-term-poss-subt-at*)
from *i* **have** *1*:($\langle$*map source ?Bs*$\rangle_\alpha$) *?x* = *source* (*?Bs*!*i*)
using *l lhs-subst-var-i* **by** (*metis length-map nth-map*)
from *i* **have** *2*:($\langle$*map source As*$\rangle_\alpha$) *?x* = *source* (*As*!*i*)
using *l lhs-subst-var-i* **by** (*metis length-map nth-map*)
from *subterm*(*5*) **show** *False* **using** *neq 1 2 p*
unfolding *Prule B source.simps source-apply-subst*[*OF to-pterm-wf-pterm*[*of lhs* α]] *source-to-pterm*
using *term-subst-eq-rev x* **by** *fastforce*
qed
from *subterm*(*1*,*2*,*3*) *co-init* **have** *As*!*i* $\perp_p$ *?Bs*!*i*
using *i* l' *orth* **unfolding** *Prule* **by** (*metis B fun-well-arg l lhs-subst-args-wf-pterm nth-mem orth supt.arg*)
}
then show *?thesis* **unfolding** *B Prule*
by (*smt* (*verit*, *best*) *case-prodI2 fst-conv in-set-zip l* l' *orthogonal.intros*(*3*) *snd-conv*)
next
case (*Prule* β *Bs*)
from *subterm*(*4*) **have** α:$\alpha = \beta$
unfolding *Prule* ‹*A* = *Prule* α *As*› *eval-term.simps* **using** *orthogonal.simps*
by (*smt* (*verit*) *Inl-inject Prule eval-term.simps*(*2*) *is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *lhs no-var-lhs.lhs-is-Fun*
no-var-lhs-axioms subterm.prems(*2*) *term.collapse*(*2*) *term.sel*(*2*) *to-pterm.simps*(*2*))
from *subterm*(*2*,*3*) **have** *l*:*length As* = *length Bs*
unfolding ‹*A* = *Prule* α *As*› *Prule* **using** α *length-args-well-Prule* **by** *blast*
{fix *i* **assume** *i*:*i* < *length As*
with *subterm*(*4*) **have** *As*!$i \cdot \sigma \perp_p$ *Bs*!$i \cdot \sigma$
unfolding ‹*A* = *Prule* α *As*› *Prule eval-term.simps* α **by** (*smt* (*verit*, *ccfv-threshold*) *Inl-inject*

Inr-not-Inl α *eval-term.simps*(2) *length-map lhs nth-map orthogonal.simps term.distinct*(1) *term.inject*(2) *to-pterm.simps*(2))
moreover from *i subterm*(5) **have** *source* (*As*!*i*) = *source* (*Bs*!*i*)
using *Prule* α ‹*A* = *Prule* α *As*› *co-init-prule subterm.prems*(1) *subterm.prems*(2) **by** *blast*
moreover from *i l subterm*(2,3) **have** *As*!*i* ∈ *wf-pterm R Bs*!*i* ∈ *wf-pterm R*
unfolding *Prule* ‹*A* = *Prule* α *As*› **by** *auto*
moreover from *i* **have** *As*!*i* ◁ *A*
unfolding ‹*A* = *Prule* α *As*› **by** *simp*
ultimately have *As*!*i* $\perp_p$ *Bs*!*i*
using *subterm*(1) **by** *simp*
}
with *l* **show** *?thesis*
unfolding α ‹*A* = *Prule* α *As*› *Prule* **by** (*simp add*: *orthogonal.intros*(2))
qed
qed
qed

end

end

5 Labels and Overlaps

theory *Labels-and-Overlaps*
imports
Orthogonal-PT
Well-Quasi-Orders.Almost-Full-Relations
begin

5.1 Labeled Proof Terms

The idea is to label function symbols in the source of a proof term that are affected by a rule symbol α with α and the distance from the root to α. Therefore, a label is a pair consisting of a rule symbol and a natural number, or it can be *None*. A labeled term is a term, where each function symbol additionally has a label associated with it.

type-synonym
('*f*, '*v*) *label* = (('*f*, '*v*) *prule* × *nat*) *option*
type-synonym
('*f*, '*v*) *term-lab* = ('*f* × ('*f*, '*v*) *label*, '*v*) *term*

fun *label-term* :: ('*f*, '*v*) *prule* ⇒ *nat* ⇒ ('*f*, '*v*) *term* ⇒ ('*f*, '*v*) *term-lab*
where
label-term α *i* (*Var x*) = *Var x*
| *label-term* α *i* (*Fun f ts*) = *Fun* (*f*, *Some* (α, *i*)) (*map* (*label-term* α (*i*+1)) *ts*)

abbreviation *labeled-lhs* :: ($'f$, $'v$) *prule* ⇒ ($'f$, $'v$) *term-lab*
  where *labeled-lhs* α ≡ *label-term* α *0* (*lhs* α)

fun *labeled-source* :: ($'f$, $'v$) *pterm* ⇒ ($'f$, $'v$) *term-lab*
  where
  labeled-source (*Var x*) = *Var x*
| *labeled-source* (*Pfun f As*) = *Fun* (*f*, *None*) (*map labeled-source As*)
| *labeled-source* (*Prule* α *As*) = (*labeled-lhs* α) · $\langle$*map labeled-source As*$\rangle_\alpha$

fun *term-lab-to-term* :: ($'f$, $'v$) *term-lab* ⇒ ($'f$, $'v$) *term*
  where
  term-lab-to-term (*Var x*) = *Var x*
| *term-lab-to-term* (*Fun f ts*) = *Fun* (*fst f*) (*map term-lab-to-term ts*)

fun *term-to-term-lab* :: ($'f$, $'v$) *term* ⇒ ($'f$, $'v$) *term-lab*
  where
  term-to-term-lab (*Var x*) = *Var x*
| *term-to-term-lab* (*Fun f ts*) = *Fun* (*f*, *None*) (*map term-to-term-lab ts*)

fun *get-label* :: ($'f$, $'v$) *term-lab* ⇒ ($'f$, $'v$) *label*
  where
  get-label (*Var x*) = *None*
| *get-label* (*Fun f ts*) = *snd f*

fun *labelposs* :: ($'f$, $'v$) *term-lab* ⇒ *pos set*
where
  labelposs (*Var x*) = {}
| *labelposs* (*Fun* (*f*, *None*) *ts*) = ($\bigcup$*i*<*length ts*. {*i* # *p* | *p*. *p* ∈ *labelposs* (*ts* ! *i*)})
| *labelposs* (*Fun* (*f*, *Some l*) *ts*) = {[]} ∪ ($\bigcup$*i*<*length ts*. {*i* # *p* | *p*. *p* ∈ *labelposs* (*ts* ! *i*)})

abbreviation *possL* :: ($'f$, $'v$) *pterm* ⇒ *pos set*
  where *possL A* ≡ *labelposs* (*labeled-source A*)

lemma *labelposs-term-to-term-lab*: *labelposs* (*term-to-term-lab t*) = {}
  by(*induct t*) *simp-all*

lemma *term-lab-to-term-lab*[*simp*]: *term-lab-to-term* (*term-to-term-lab t*) = *t*
proof(*induct t*)
  case (*Fun f ts*)
  then show *?case*
    unfolding *term-lab-to-term.simps term-to-term-lab.simps fst-conv* **by** (*simp add*: *map-nth-eq-conv*)
qed *simp*

lemma *term-lab-to-term-subt-at*:
  assumes *p* ∈ *poss t*
  shows *term-lab-to-term t* |-*p* = *term-lab-to-term* (*t*|-*p*)
  using *assms* **proof**(*induct p arbitrary*:*t*)

case (*Cons i p*)
from *args-poss*[*OF Cons*(*2*)] **obtain** *f ts* **where** *f*:*t* = *Fun f ts* **and**
p:*p* ∈ *poss* (*ts* ! *i*) **and** *i*:*i* < *length ts* **by** *blast*
from *Cons*(*1*)[*OF p*] *i* **show** *?case*
unfolding *f term-lab-to-term.simps* **by** *simp*
qed *simp*

lemma *vars-term-labeled-lhs*: *vars-term* (*label-term* α *i t*) = *vars-term t*
by (*induct t arbitrary*:*i*) *simp-all*

lemma *vars-term-list-labeled-lhs*: *vars-term-list* (*label-term* α *i t*) = *vars-term-list*
t
proof (*induct t arbitrary*:*i*)
case (*Fun f ts*)
show *?case* **unfolding** *label-term.simps vars-term-list.simps* **using** *Fun*
by (*metis* (*mono-tags*, *lifting*) *length-map map-nth-eq-conv nth-mem*)
qed (*simp add*: *vars-term-list.simps*(*1*))

lemma *var-poss-list-labeled-lhs*: *var-poss-list* (*label-term* α *i t*) = *var-poss-list t*
proof (*induct t arbitrary*:*i*)
case (*Fun f ts*)
then have *ts*:*map* (*var-poss-list* ∘ *label-term* α (*i* + *1*)) *ts* = *map var-poss-list*
ts
by *auto*
then show *?case*
unfolding *label-term.simps var-poss-list.simps map-map ts* **by** *simp*
qed (*simp add*: *poss-list.simps*(*1*))

lemma *var-labeled-lhs*[*simp*]: *vars-distinct* (*label-term* α *i t*) = *vars-distinct t*
by (*simp add*: *vars-term-list-labeled-lhs*)

lemma *labelposs-subt-at*:
assumes *q* ∈ *poss t p* ∈ *labelposs* (*t*|-*q*)
shows *q*@*p* ∈ *labelposs t*
using *assms* **proof**(*induct t arbitrary*:*q*)
case (*Fun f ts*)
then show *?case* **proof**(*cases q*)
case (*Cons i q'*)
with *Fun*(*2*) **have** *i*:*i* < *length ts* **and** *q'*:*q'* ∈ *poss* (*ts*!*i*)
by *simp*+
with *Fun*(*3*) **have** *p* ∈ *labelposs* ((*ts*!*i*)|-*q'*)
unfolding *Cons* **by** *simp*
with *Fun*(*1*) *i q'* **have** *IH*:*q'*@*p* ∈ *labelposs* (*ts*!*i*)
using *nth-mem* **by** *blast*
obtain *f' lab* **where** *f*:*f* = (*f'*, *lab*)
by *fastforce*
then show *?thesis* **proof**(*cases lab*)
case *None*
show *?thesis*

unfolding *f Cons None labelposs.simps* **using** *i IH* **by** *simp*
next
case (*Some a*)
then show *?thesis*
unfolding *f Cons Some labelposs.simps* **using** *i IH* **by** *simp*
qed
qed *simp*
qed *simp*

lemma *var-label-term*:
assumes $p \in$ *poss t* **and** *t*|-*p* = *Var x*
shows *label-term* α *n t* |-*p* = *Var x*
using *assms* **proof**(*induct t arbitrary*:*p n*)
case (*Fun f ts*)
then obtain *i p'* **where** *p'*:*i* < *length ts p* = *i*#*p' p'* $\in$ *poss* (*ts*!*i*)
by *auto*
then show *?case*
unfolding *label-term.simps p'*(*2*) *subt-at.simps* **using** *Fun*(*1*,*3*) *p'*(*2*) **by** *force*
qed *simp*

lemma *get-label-label-term*:
assumes $p \in$ *fun-poss t*
shows *get-label* (*label-term* α *n t*|-*p*) = *Some* (α, *n* + *size p*)
using *assms* **proof**(*induct t arbitrary*: *n p*)
case (*Fun f ts*)
show *?case* **proof**(*cases p*)
case (*Cons i p'*)
with *Fun*(*2*) **have** *i*:*i* < *length ts* **and** *p'*:*p'* $\in$ *fun-poss* (*ts*!*i*) **by** *simp*+
with *Fun*(*1*) **have** *get-label* (*label-term* α (*n*+*1*) (*ts*!*i*) |- *p'*) = *Some* (α, *n* + *1* + *size p'*) **by** *simp*
then show *?thesis* **unfolding** *Cons label-term.simps subt-at.simps* **using** *i* **by** *auto*
qed *simp*
qed *simp*

lemma *linear-label-term*:
assumes *linear-term t*
shows *linear-term* (*label-term* α *n t*)
using *assms* **proof**(*induct t arbitrary*:*n*)
case (*Fun f ts*)
from *Fun*(*2*) **have** (*is-partition* (*map vars-term ts*))
by *simp*
then have *is-partition* (*map vars-term* (*map* (*label-term* α (*Suc n*)) *ts*))
by (*metis* (*mono-tags*, *lifting*) *length-map map-nth-eq-conv vars-term-labeled-lhs*)
moreover {**fix** *t* **assume** *t*:*t* $\in$ *set ts*
with *Fun*(*2*) **have** *linear-term t*
by *simp*
with *Fun*(*1*) **have** *linear-term* (*label-term* α (*Suc n*) *t*)
using *t* **by** *blast*

}
ultimately show *?case* unfolding *label-term.simps* by *simp*
qed *simp*

lemma *var-term-lab-to-term*:
assumes $p \in poss\ t$ and $t|\text{-}p = Var\ x$
shows *term-lab-to-term* $t\ |\text{-}p = Var\ x$
using *assms* proof(*induct t arbitrary*:*p*)
case (*Fun f ts*)
then obtain $i\ p'$ where p':$i < length\ ts\ p = i\#p'\ p' \in poss\ (ts!i)$
by *auto*
then show *?case*
unfolding *term-lab-to-term.simps* $p'(2)$ *subt-at.simps* using *Fun*(*1*,*3*) $p'(2)$
by *force*
qed *simp*

lemma *poss-term-lab-to-term*[*simp*]: *poss t* = *poss* (*term-lab-to-term t*)
by(*induct t*) *auto*

lemma *fun-poss-term-lab-to-term*[*simp*]: *fun-poss t* = *fun-poss* (*term-lab-to-term t*)
by(*induct t*) *auto*

lemma *vars-term-list-term-lab-to-term*: *vars-term-list t* = *vars-term-list* (*term-lab-to-term t*)
proof(*induct t*)
case (*Var x*)
then show *?case*
by (*simp add*: *vars-term-list.simps*(*1*))
next
case (*Fun f ts*)
then show *?case* unfolding *vars-term-list.simps term-lab-to-term.simps*
by (*smt* (*verit*, *best*) *length-map map-eq-conv′ nth-map nth-mem*)
qed

lemma *vars-term-list-term-to-term-lab*: *vars-term-list* (*term-to-term-lab t*) = *vars-term-list t*
proof(*induct t*)
case (*Var x*)
then show *?case*
by (*simp add*: *vars-term-list.simps*(*1*))
next
case (*Fun f ts*)
then show *?case* unfolding *vars-term-list.simps term-to-term-lab.simps*
by (*metis* (*mono-tags*, *lifting*) *length-map map-nth-eq-conv nth-mem*)
qed

lemma *linear-term-to-term-lab*:
assumes *linear-term t*

shows *linear-term* (*term-to-term-lab t*)
using *assms* **proof**(*induct t*)
case (*Fun f ts*)
then show *?case* **unfolding** *term-to-term-lab.simps linear-term.simps*
by (*smt* (*verit*, *best*) *imageE length-map list.set-map map-nth-eq-conv set-vars-term-list vars-term-list-term-to-term-lab*)
qed *simp*

lemma *var-poss-list-term-lab-to-term*: *var-poss-list t* = *var-poss-list* (*term-lab-to-term t*)
proof(*induct t*)
case (*Var x*)
then show *?case*
by (*simp add*: *poss-list.simps*(*1*))
next
case (*Fun f ts*)
then have ∗:(*map var-poss-list ts*) = (*map var-poss-list* (*map term-lab-to-term ts*))
by *auto*
then show *?case* **unfolding** *term-lab-to-term.simps var-poss-list.simps length-map* ∗
by *blast*
qed

lemma *label-poss-labeled-lhs*:
assumes $p \in$ *fun-poss* (*label-term* α *n t*)
shows $p \in$ *labelposs* (*label-term* α *n t*)
using *assms* **proof**(*induct t arbitrary*:*p n*)
case (*Fun f ts*)
then show *?case* **proof**(*cases p*)
case (*Cons i p′*)
from *Fun*(*2*) **have** *i*:$i <$ *length ts*
unfolding *Cons* **by** *simp*
with *Fun*(*2*) **have** $p' \in$ *fun-poss* (*label-term* α (*n+1*) (*ts*!*i*))
unfolding *Cons* **by** *auto*
with *i* **have** $p' \in$ *labelposs* (*label-term* α (*n+1*) (*ts*!*i*))
using *Fun*(*1*) **by** *simp*
with *i* **show** *?thesis*
unfolding *Cons label-term.simps labelposs.simps* **by** *simp*
qed *simp*
qed *simp*

lemma *labeled-var*:
assumes *source A* = *Var x*
shows *labeled-source A* = *Var x*
using *assms* **proof**(*induct A*)
case (*Prule* α *As*)
then show *?case* **proof**(*cases As* = [])
case *True*

from *Prule*(*2*) **have** *lhs* α = *Var x*
unfolding *source.simps True list.map* **by** *simp*
with *True* **show** *?thesis*
by *simp*
next
case *False*
then obtain *a as* **where** *as*:*As* = *a* # *as*
using *list.exhaust* **by** *blast*
from *Prule*(*2*) **obtain** *y* **where** *y*:*lhs* α = *Var y*
using *is-Var-def* **by** *fastforce*
from *Prule*(*2*) **have** *source a* = *Var x*
unfolding *source.simps y as single-var* **by** *simp*
with *Prule*(*1*) *as* **have** *labeled-source a* = *Var x*
by *simp*
then show *?thesis*
unfolding *labeled-source.simps as y single-var* **by** *simp*
qed
qed *simp-all*

lemma *labelposs-subs-fun-poss*: *labelposs t* $\subseteq$ *fun-poss t*
proof(*induct t*)
case (*Fun fl ts*)
then show *?case* **proof**(*cases snd fl*)
case *None*
then obtain *f* **where** *f*:*fl* = (*f*, *None*)
by (*metis prod.collapse*)
then have *labelposs* (*Fun fl ts*) = ($\bigcup$*i*<*length ts*. {*i* # *p* |*p*. *p* $\in$ *labelposs* (*ts* ! *i*)})
by *simp*
also have ... $\subseteq$ ($\bigcup$*i*<*length ts*. {*i* # *p* |*p*. *p* $\in$ *fun-poss* (*ts* ! *i*)}) **using** *Fun*
by (*smt SUP-mono basic-trans-rules*(*31*) *lessThan-iff mem-Collect-eq nth-mem subsetI*)
finally show *?thesis*
by *auto*
next
case (*Some l*)
then obtain *f* **where** *f*:*fl* = (*f*, *Some l*)
by (*metis prod.collapse*)
then have *labelposs* (*Fun fl ts*) = {[]} $\cup$ ($\bigcup$*i*<*length ts*. {*i* # *p* |*p*. *p* $\in$ *labelposs* (*ts* ! *i*)})
by *simp*
also have ... $\subseteq$ {[]} $\cup$ ($\bigcup$*i*<*length ts*. {*i* # *p* |*p*. *p* $\in$ *fun-poss* (*ts* ! *i*)}) **using** *Fun*
by (*smt SUP-mono basic-trans-rules*(*31*) *lessThan-iff mem-Collect-eq nth-mem subsetI sup-mono*)
finally show *?thesis*
by *auto*
qed
qed *simp*

```
lemma labelposs-subs-poss[simp]: labelposs t ⊆ poss t
  using labelposs-subs-fun-poss fun-poss-imp-poss by blast

lemma get-label-imp-labelposs:
  assumes p ∈ poss t and get-label (t|-p) ≠ None
  shows p ∈ labelposs t
  using assms proof(induct p arbitrary:t)
  case Nil
  then show ?case unfolding subt-at.simps
   by (smt UnCI get-label.elims insert-iff labelposs.elims prod.sel(2) term.distinct(1)
term.inject(2))
next
  case (Cons i p)
  then obtain f ts where t:t = Fun f ts and p ∈ poss (ts ! i) and i:i < length ts
    using args-poss by blast
  with Cons(1,3) have p ∈ labelposs (ts!i)
    by simp
  with i have p:i # p ∈ (⋃i<length ts. {i # p |p. p ∈ labelposs (ts ! i)})
    by blast
  then show ?case proof(cases snd f)
    case None
    with p show ?thesis unfolding t using labelposs.simps(2)
      by (metis (mono-tags, lifting) prod.collapse)
  next
    case (Some a)
    with p show ?thesis unfolding t using labelposs.simps(3)
      by (smt UN-iff Un-iff mem-Collect-eq prod.collapse)
  qed
qed

lemma labelposs-obtain-label:
  assumes p ∈ labelposs t
  shows ∃α m. get-label (t|-p) = Some(α, m)
  using assms proof(induct t arbitrary: p)
  case (Fun fl ts)
  then show ?case proof(cases p)
    case Nil
    {fix f assume f:fl = (f, None)
      from Fun(2) have False unfolding Nil f labelposs.simps(2)
        by blast
    }
    with Nil show ?thesis
      by (metis eq-snd-iff get-label.simps(2) option.exhaust subt-at.simps(1))
  next
    case (Cons i q)
    with Fun(2) have iq:i # q ∈ labelposs (Fun fl ts)
      by simp
    then have i:i < length ts
```

using *labelposs-subs-poss* **by** *fastforce*
with *iq* **have** $i \# q \in \{i \# p \mid p.\ p \in labelposs\ (ts\ !\ i)\}$ **proof**(*cases snd fl*)
case (*Some a*)
then obtain $f\ \alpha\ n$ **where** $f{:}fl = (f,\ Some\ (\alpha,\ n))$
by (*metis eq-snd-iff*)
from *iq* **show** *?thesis* **unfolding** *f labelposs.simps*
by *blast*
qed (*smt UN-iff labelposs.simps(2) list.inject mem-Collect-eq prod.collapse*)
with *i Fun(1) Cons* **show** *?thesis*
by *simp*
qed
qed *simp*

lemma *possL-obtain-label*:
assumes $p \in possL\ A$
shows $\exists \alpha\ m.\ get\text{-}label\ ((labeled\text{-}source\ A)|\text{-}p) = Some(\alpha,\ m)$
using *assms labelposs-obtain-label* **by** *blast*

lemma *labeled-source-apply-subst*:
assumes $A \in wf\text{-}pterm\ R$
shows $labeled\text{-}source\ (A \cdot \sigma) = (labeled\text{-}source\ A) \cdot (labeled\text{-}source \circ \sigma)$
using *assms* **proof**(*induct A*)
case ($3\ \alpha\ As$)
have $id{:}\forall\ x \in vars\text{-}term\ (labeled\text{-}lhs\ \alpha).\ ((\langle map\ (labeled\text{-}source \circ (\lambda t.\ t \cdot \sigma))\ As\rangle_\alpha)$ $x = (\langle map\ labeled\text{-}source\ As\rangle_\alpha \circ_s (labeled\text{-}source \circ \sigma))\ x$
proof−
have $vars{:}vars\text{-}term\ (labeled\text{-}lhs\ \alpha) = set\ (var\text{-}rule\ \alpha)$ **using** *vars-term-labeled-lhs* **by** *simp*
{ fix *i* **assume** $i{:}i < length\ (var\text{-}rule\ \alpha)$
with *3* **have** $(\langle map\ (labeled\text{-}source \circ (\lambda t.\ t \cdot \sigma))\ As\rangle_\alpha)\ ((var\text{-}rule\ \alpha)!i) = labeled\text{-}source\ ((As!i) \cdot \sigma)$
by (*simp add: mk-subst-distinct*)
also have ... $= labeled\text{-}source\ (As!i) \cdot (labeled\text{-}source \circ \sigma)$
using *3 i* **by** (*metis nth-mem*)
also have ... $= (\langle map\ labeled\text{-}source\ As\rangle_\alpha \circ_s (labeled\text{-}source \circ \sigma))\ ((var\text{-}rule\ \alpha)!i)$
using *3 i* **unfolding** *subst-compose-def* **by** (*simp add: mk-subst-distinct*)
finally have $(\langle map\ (labeled\text{-}source \circ (\lambda t.\ t \cdot \sigma))\ As\rangle_\alpha)\ ((var\text{-}rule\ \alpha)!i) = (\langle map\ labeled\text{-}source\ As\rangle_\alpha \circ_s (labeled\text{-}source \circ \sigma))\ ((var\text{-}rule\ \alpha)!i)$ **.**
} with *vars* **show** *?thesis* **by** (*metis in-set-idx*)
qed
have $labeled\text{-}source\ ((Prule\ \alpha\ As) \cdot \sigma) = (labeled\text{-}lhs\ \alpha) \cdot \langle map\ (labeled\text{-}source \circ (\lambda t.\ t \cdot \sigma))\ As\rangle_\alpha$
unfolding *eval-term.simps(2)* **by** *simp*
also have ... $= (labeled\text{-}lhs\ \alpha) \cdot (\langle map\ labeled\text{-}source\ As\rangle_\alpha \circ_s (labeled\text{-}source \circ \sigma))$
using *id* **by** (*meson term-subst-eq*)
also have ... $= (labeled\text{-}source\ (Prule\ \alpha\ As)) \cdot (labeled\text{-}source \circ \sigma)$ **by** *simp*
finally show *?case* **.**

qed *simp-all*

lemma *labelposs-apply-subst*:
labelposs $(s \cdot \sigma)$ = *labelposs s* $\cup$ $\{p@q|$ *p q x. p* $\in$ *var-poss s* $\wedge$ *s*$|$-*p* = *Var x* $\wedge$ *q* $\in$ *labelposs* $(\sigma\ x)\}$
proof(*induct s*)
case (*Fun f ts*)
obtain f' *l* **where** *f*:*f* = (f', l) **by** *fastforce*
let *?lp1*=$\bigcup$ *i*<*length ts*. $\{i\ \#\ p\ |p.\ p \in$ *labelposs* (*ts* ! *i*)$\}$
let *?lp2*=$\bigcup$ *i*<*length ts*. $\{i\#(p@q)|$ *p q x. p* $\in$ *var-poss* (*ts*!*i*) $\wedge$ (*ts*!*i*)$|$-*p* = *Var x* $\wedge$ *q* $\in$ *labelposs* $(\sigma\ x)\}$
$\{$**fix** *i* **assume** *i*:*i* < *length ts*
with *Fun* **have** $\{i\ \#\ p\ |p.\ p \in$ *labelposs* (*ts* ! *i* $\cdot\ \sigma$)$\}$ = $\{i\ \#\ p|\ p.\ p \in$ *labelposs* (*ts*!*i*) $\cup$ $\{p@q|$ *p q x. p* $\in$ *var-poss* (*ts*!*i*) $\wedge$ (*ts*!*i*)$|$-*p* = *Var x* $\wedge$ *q* $\in$ *labelposs* $(\sigma\ x)\}\}$
by *auto*
then have $\{i\ \#\ p\ |p.\ p \in$ *labelposs* (*map* ($\lambda s.\ s \cdot \sigma$) *ts* ! *i*)$\}$ = $\{i\ \#\ p|\ p.\ p \in$ *labelposs* (*ts*!*i*)$\}$ $\cup$ $\{i\#(p@q)|$ *p q x. p* $\in$ *var-poss* (*ts*!*i*) $\wedge$ (*ts*!*i*)$|$-*p* = *Var x* $\wedge$ *q* $\in$ *labelposs* $(\sigma\ x)\}$
unfolding *Un-iff* **using** *i* **by** *auto*
$\}$**note** *IH*=*this*
$\{$**fix** *i* **assume** *i*:*i* < *length ts*
let *?l*=$\{i\#(p@q)|$ *p q x. p* $\in$ *var-poss* (*ts*!*i*) $\wedge$ (*ts*!*i*)$|$-*p* = *Var x* $\wedge$ *q* $\in$ *labelposs* $(\sigma\ x)\}$
let *?r*=$\{p@q|$ *p q x. p* $\in$ $\{i\ \#\ p\ |p.\ p \in$ *var-poss* (*ts* ! *i*)$\}$ $\wedge$ (*Fun f ts*)$|$-*p* = *Var x* $\wedge$ *q* $\in$ *labelposs* $(\sigma\ x)\}$
have *?l* = *?r* **proof**
show *?l* $\subseteq$ *?r*
by (*smt* (*verit*, *ccfv-SIG*) *Collect-mono-iff Cons-eq-appendI mem-Collect-eq subt-at.simps*(*2*))
show *?r* $\subseteq$ *?l*
by (*smt* (*verit*, *best*) *Collect-mono-iff Cons-eq-appendI mem-Collect-eq subt-at.simps*(*2*))
qed
$\}$
then have *lp2*:$\{p@q|$ *p q x. p* $\in$ *var-poss* (*Fun f ts*) $\wedge$ (*Fun f ts*)$|$-*p* = *Var x* $\wedge$ *q* $\in$ *labelposs* $(\sigma\ x)\}$ = *?lp2*
unfolding *var-poss.simps* **by** *auto*
show *?case* **proof**(*cases l*)
case *None*
have *labelposs* (*Fun f ts* $\cdot\ \sigma$) = *?lp1* $\cup$ *?lp2*
unfolding *eval-term.simps f None labelposs.simps length-map* **using** *IH* **by** *auto*
moreover have *labelposs* (*Fun f ts*) = *?lp1*
unfolding *f None* **by** *simp*
ultimately show *?thesis* **using** *lp2* **by** *simp*
next
case (*Some a*)
have *labelposs* (*Fun f ts* $\cdot\ \sigma$) = $\{[]\}$ $\cup$ *?lp1* $\cup$ *?lp2*
unfolding *eval-term.simps f Some labelposs.simps length-map* **using** *IH* **by**

auto
  moreover have *labelposs* (*Fun f ts*) = {[]} ∪ *?lp1*
    unfolding *f Some* **by** *simp*
  ultimately show *?thesis* **using** *lp2* **by** *simp*
 qed
qed *simp*

lemma *possL-apply-subst*:
 assumes $A \cdot \sigma \in$ *wf-pterm R*
 shows *possL* $(A \cdot \sigma)$ = *possL A* ∪ {*p*@*q*| *p q x*. *p* ∈ *var-poss* (*labeled-source A*) ∧ (*labeled-source A*)|-*p* = *Var x* ∧ *q* ∈ *possL* (σ *x*)}
proof−
 from *assms* **have** ∗:*labeled-source* $(A \cdot \sigma)$ = *labeled-source A* · (*labeled-source* ∘ σ)
  using *labeled-source-apply-subst subst-imp-well-def* **by** *blast*
 then show *?thesis* **unfolding** ∗ *labelposs-apply-subst*
  by *auto*
qed

lemma *label-term-to-term*[*simp*]: *term-lab-to-term* (*label-term* α *n t*) = *t*
 by(*induct t arbitrary*:α *n*)(*simp-all add*: *map-nth-eq-conv*)

lemma *fun-poss-label-term*: *p* ∈ *fun-poss* (*label-term* β *n t*) ⟷ *p* ∈ *fun-poss t*
proof
 {**assume** *p* ∈ *fun-poss* (*label-term* β *n t*)
  then show *p* ∈ *fun-poss t* **proof**(*induct t arbitrary*:*n p*)
   case (*Fun f ts*)
   then show *?case* **by**(*cases p*) *auto*
  qed *simp*
 }
 {**assume** *p* ∈ *fun-poss t*
  then show *p* ∈ *fun-poss* (*label-term* β *n t*) **proof**(*induct t arbitrary*:*n p*)
   case (*Fun f ts*)
   then show *?case* **by**(*cases p*) *auto*
  qed *simp*
 }
qed

lemma *term-lab-to-term-subst*: *term-lab-to-term* $(t \cdot \sigma)$ = *term-lab-to-term t* · (*term-lab-to-term* ∘ σ)
proof(*induct t*)
 case (*Fun f As*)
 then show *?case* **unfolding** *eval-term.simps*(*2*) *term-lab-to-term.simps*
  by *fastforce*
qed *simp*

lemma *labeled-source-to-term*[*simp*]: *term-lab-to-term* (*labeled-source A*) = *source A*
proof(*induct A*)

case (*Prule* α *As*)
have *term-lab-to-term* $\circ$ $\langle$*map labeled-source As*$\rangle_\alpha$ = $\langle$*map* (*term-lab-to-term* $\circ$ *labeled-source*) *As*$\rangle_\alpha$
by *simp*
also have ... = $\langle$*map source As*$\rangle_\alpha$ **using** *Prule*
by (*metis* (*mono-tags*, *lifting*) *comp-apply map-eq-conv*)
finally show *?case* **unfolding** *labeled-source.simps source.simps*
by (*simp add*: *term-lab-to-term-subst*)
qed *simp-all*

lemma *possL-subset-poss-source*: *possL A* $\subseteq$ *poss* (*source A*)
using *poss-term-lab-to-term labeled-source-to-term labelposs-subs-poss*
by *metis*

lemma *labeled-source-pos*:
assumes *p* $\in$ *poss s* **and** *term-lab-to-term t* = *s*
shows *term-lab-to-term* (*t*|-*p*) = *s*|-*p*
using *assms* **proof**(*induct p arbitrary*:*s t*)
case (*Cons i p*)
from *Cons*(*2*) **obtain** *f ss* **where** *s*:*s* = *Fun f ss*
using *args-poss* **by** *blast*
with *Cons*(*2*) **have** *p*:*p* $\in$ *poss* (*ss*!*i*)
by *force*
from *Cons*(*3*) *s* **obtain** *label ts* **where** *t*:*t* = *Fun* (*f*, *label*) *ts*
by (*metis args-poss local.Cons*(*2*) *poss-term-lab-to-term prod.collapse term.inject*(*2*) *term-lab-to-term.simps*(*2*))
with *Cons*(*2*,*3*) *s* **have** *term-lab-to-term* (*ts*!*i*) = *ss*!*i*
by *auto*
with *Cons*(*1*) *p* **show** *?case* **unfolding** *s t*
by *simp*
qed *simp*

lemma *get-label-at-fun-poss-subst*:
assumes *p* $\in$ *fun-poss t*
shows *get-label* ((*t* $\cdot$ σ)|-*p*) = *get-label* (*t*|-*p*)
using *assms fun-poss-fun-conv fun-poss-imp-poss* **by** *fastforce*

lemma *labeled-source-simple-pterm*:*possL* (*to-pterm t*) = {}
by(*induct t*) *simp-all*

lemma *label-term-increase*:
assumes *s* = (*label-term* α *n t*) $\cdot$ σ **and** *p* $\in$ *fun-poss t*
shows *get-label* (*s*|-*p*) = *Some* (α, *n* + *length p*)
using *assms* **proof**(*induct p arbitrary*: *s t n*)
case *Nil*
then obtain *f ts* **where** *t* = *Fun f ts*
by (*metis fun-poss-list.simps*(*1*) *in-set-simps*(*3*) *is-FunE is-Var-def set-fun-poss-list*)
with *Nil*(*1*) **show** *?case*
by *simp*

next
  case (*Cons i p*)
  then obtain *f ts* **where** *f*:*t* = *Fun f ts* **and** *i*:*i*<*length ts*
    by (*meson args-poss fun-poss-imp-poss*)
  with *Cons*(*3*) **have** *p*:*p* ∈ *fun-poss* (*ts*!*i*)
    by *auto*
  let *?s′* = (*label-term* α (*n*+*1*) (*ts*!*i*)) · σ
  from *Cons*(*1*) *p* **have** *get-label* (*?s′*|- *p*) = *Some* (α, *n* + *1* + *length p*)
    by *blast*
  with *i* **show** *?case* **unfolding** *Cons*(*2*) *f*
    by *simp*
qed

The number attached to a labeled function symbol cannot exceed the depth of that function symbol.

lemma *label-term-max-value*:
 assumes *p* ∈ *poss* (*labeled-source A*) **and** *get-label* ((*labeled-source A*)|-*p*) = *Some* (α, *n*)
    and *A* ∈ *wf-pterm R*
  shows *n* ≤ *length p*
  using *assms* **proof**(*induct A arbitrary*: *p*)
  case (*Pfun f As*)
  then show *?case* **proof**(*cases p*)
    case (*Cons i q*)
    with *Pfun*(*2*) **have** *i*:*i* < *length As* **by** *simp*
    with *Pfun*(*3*) **have** *lab*:*get-label* (*labeled-source* (*As*!*i*) |- *q*) = *Some* (α, *n*)
      unfolding *Cons* **by** *simp*
    with *Pfun*(*2*) *i* **have** *q* ∈ *poss* (*labeled-source* (*As*!*i*))
      unfolding *Cons* **by** *auto*
    with *Pfun*(*1*,*4*) *Cons i lab* **show** *?thesis*
      using *nth-mem fun-well-arg* **by** *fastforce*
  qed *simp*
next
  case (*Prule* β *As*)
  from *Prule*(*2*) **consider** *p* ∈ *fun-poss* (*labeled-lhs* β) | (∃ *p1 p2 x*. *p* = *p1*@*p2*
                                   ∧ *p1* ∈ *poss* (*labeled-lhs* β) ∧ (*labeled-lhs* β)|-*p1* = *Var x*
                                   ∧ *p2* ∈ *poss* ((⟨*map labeled-source As*⟩$_\beta$) *x*)
                                   ∧ (*labeled-source* (*Prule* β *As*))|-*p* = ((⟨*map labeled-source As*⟩$_\beta$) *x*)|-*p2*)
   unfolding *labeled-source.simps* **by** (*meson poss-is-Fun-fun-poss poss-subst-choice*)
  then show *?case* **proof**(*cases*)
    case *1*
    then have *p* ∈ *fun-poss* (*lhs* β)
      by (*simp add*: *fun-poss-label-term*)
    then have *get-label* ((*labeled-source* (*Prule* β *As*))|-*p*) = *Some* (β, *length p*)
      unfolding *labeled-source.simps* **by** (*simp add*: *label-term-increase*)
    with *Prule*(*3*) **show** *?thesis* **by** *auto*
  next

case *2*
then obtain *p1 p2 x* **where** *p1p2*:*p* = *p1* @ *p2* **and** *x*:*p1* ∈ *poss* (*labeled-lhs* β) ∧ *labeled-lhs* β |- *p1* = *Var x*
and *p2*:*p2* ∈ *poss* ((⟨*map labeled-source As*⟩$_\beta$) *x*)
and *lab*:*labeled-source* (*Prule* β *As*) |- *p* = (⟨*map labeled-source As*⟩$_\beta$) *x* |- *p2*
by *blast*
from *Prule*(*4*) **have** *l*:*length As* = *length* (*var-rule* β)
using *wf-pterm.simps* **by** *fastforce*
from *x* **have** *x* ∈ *vars-term* (*lhs* β)
by (*metis subt-at-imp-supteq subteq-Var-imp-in-vars-term vars-term-labeled-lhs*)
with *x* **obtain** *i* **where** *i*:*i* < *length* (*var-rule* β) ∧ (*var-rule* β)!*i* = *x*
by (*metis in-set-conv-nth set-vars-term-list vars-term-list-vars-distinct*)
with *l* **have** ∗:(⟨*map labeled-source As*⟩$_\beta$) *x* = *labeled-source* (*As*!*i*)
by (*metis* (*no-types*, *lifting*) *length-map lhs-subst-var-i nth-map*)
with *Prule*(*3*) *lab* **have** *get-label* ((*labeled-source* (*As*!*i*))|-*p2*) = *Some* (α, *n*)
by *simp*
with *Prule*(*1*,*4*) *p2* ∗ *i l* **have** *n* ≤ *length p2*
by (*metis fun-well-arg nth-mem*)
with ∗ *p1p2 lab i l* **show** *?thesis* **by** *force*
qed
qed *simp*

The labels decrease when moving up towards the root from a labeled function symbol.

lemma *label-decrease*:
assumes *p*@*q* ∈ *poss* (*labeled-source A*)
and *get-label* ((*labeled-source A*)|-(*p*@*q*)) = *Some* (α, *length q* + *n*)
and *A* ∈ *wf-pterm R*
shows *get-label* ((*labeled-source A*)|-*p*) = *Some* (α, *n*)
using *assms* **proof**(*induct A arbitrary*: *p q*)
case (*Pfun f As*)
then show *?case* **proof**(*cases p*)
case *Nil*
from *Pfun*(*2*,*3*) **obtain** *i q′* **where** *iq′*:*q* = *i*#*q′* **and** *i*:*i* < *length As* **and** *q′*:*q′* ∈ *poss* (*labeled-source* (*As*!*i*))
unfolding *Nil* **by** *auto*
with *Pfun*(*2*,*3*) **have** *get-label* (*labeled-source* (*As*!*i*) |- (*q′*)) = *Some* (α, *length q* + *n*)
unfolding *Nil* **by** *auto*
with *iq′ q′* **have** *False*
using *label-term-max-value Pfun*(*4*) *i fun-well-arg* **by** (*metis le-imp-less-Suc length-nth-simps*(*2*) *not-add-less1 nth-mem*)
then show *?thesis* **by** *simp*
next
case (*Cons i p′*)
with *Pfun*(*2*) **have** *ip′*:*p* = *i*#*p′* **and** *i*:*i* < *length As*
by *auto*
with *Pfun*(*2*) **have** *p′*:*p′*@*q* ∈ *poss* (*labeled-source* (*As*!*i*))
by *simp*

from *Pfun(3) i ip'* **have** *get-label (labeled-source (As!i) |- (p'@q))* = *Some* (α, *length q* + *n*)
by *simp*
with *Pfun(1,4) p' i* **have** *get-label ((labeled-source (As!i))|-p')* = *Some* (α, *n*)
by (*metis fun-well-arg nth-mem*)
then show *?thesis*
using *i ip'* **by** *fastforce*
qed
next
case (*Prule* β *As*)
from *Prule(2)* **consider** *p@q* $\in$ *fun-poss* (*labeled-lhs* β) | ($\exists$ *p1 p2 x. p@q* = *p1@p2*
$\wedge$ *p1* $\in$ *poss* (*labeled-lhs* β) $\wedge$ (*labeled-lhs* β)|-*p1* = *Var x*
$\wedge$ *p2* $\in$ *poss* (($\langle$*map labeled-source As*$\rangle_\beta$) *x*)
$\wedge$ (*labeled-source* (*Prule* β *As*))|-(*p@q*) = (($\langle$*map labeled-source As*$\rangle_\beta$) *x*)|-*p2*)
unfolding *labeled-source.simps* **by** (*meson poss-is-Fun-fun-poss poss-subst-choice*)
then show *?case* **proof**(*cases*)
case *1*
then have *lab:get-label* ((*labeled-source* (*Prule* β *As*))|-(*p@q*)) = *Some* (β, *length p* + *length q*)
by (*simp add: fun-poss-label-term label-term-increase*)
from *1* **have** *p* $\in$ *fun-poss* (*labeled-lhs* β) **proof**(*cases q*)
case (*Cons a list*)
then show *?thesis*
using *1 fun-poss-append-poss fun-poss-imp-poss* **by** *blast*
qed *simp*
with *Prule(3) lab* **show** *?thesis*
by (*simp add: fun-poss-label-term label-term-increase*)
next
case *2*
then obtain *p1 p2 x* **where** *p1p2:p@q* = *p1* @ *p2* **and** *x:p1* $\in$ *poss* (*labeled-lhs* β) $\wedge$ *labeled-lhs* β |- *p1* = *Var x*
and *p2:p2* $\in$ *poss* (($\langle$*map labeled-source As*$\rangle_\beta$) *x*)
and *lab:labeled-source* (*Prule* β *As*) |-(*p@q*) = ($\langle$*map labeled-source As*$\rangle_\beta$) *x* |- *p2*
by *blast*
from *Prule(4)* **have** *l:length As* = *length* (*var-rule* β)
using *wf-pterm.simps* **by** *fastforce*
from *x* **have** *x* $\in$ *vars-term* (*lhs* β)
by (*metis subt-at-imp-supteq subteq-Var-imp-in-vars-term vars-term-labeled-lhs*)
then obtain *i* **where** *i:i* < *length* (*var-rule* β) $\wedge$ (*var-rule* β)!*i* = *x*
by (*metis in-set-idx set-vars-term-list vars-term-list-vars-distinct*)
with *l* **have** $*$:($\langle$*map labeled-source As*$\rangle_\beta$) *x* = *labeled-source* (*As!i*)
by (*metis* (*no-types, lifting*) *length-map lhs-subst-var-i nth-map*)
with *Prule(3) lab* **have** *as-i:get-label* ((*labeled-source* (*As!i*))|-*p2*) = *Some* (α, *length q* + *n*)
by *simp*

have *p1-above-p*:*p1* $\leq_p$ *p* **proof**(*rule ccontr*)
assume ¬ *p1* $\leq_p$ *p*
with *p1p2* **have** *length p* < *length p1*
by (*metis less-eq-pos-simps*(*1*) *pos-cases pos-less-eq-append-not-parallel prefix-smaller*)
with *p1p2* **have** *le*:*length p2* < *length q*
using *length-append* **by** (*metis add.commute less-add-eq-less*)
with *as-i Prule*(*4*) ∗ *i l p2* **have** *length q* + *n* ≤ *length p2*
by (*metis fun-well-arg label-term-max-value nth-mem*)
with *le* **show** *False* **by** *linarith*
qed
let *?p′*=*the* (*remove-prefix p1 p*)
from *p1-above-p p1p2* **have** *p2′*:*p2* = *?p′* @ *q*
by (*metis append-assoc pos-diff-def prefix-pos-diff same-append-eq*)
then have *lab′*:*labeled-source* (*Prule* β *As*) |-(*p1*@*?p′*) = (⟨*map labeled-source As*⟩$_\beta$) *x* |-*?p′*
using *x p1p2 Prule*(*2*) **unfolding** *labeled-source.simps*
by (*metis* (*mono-tags*, *lifting*) *labeled-source.simps*(*3*) *poss-append-poss eval-term.simps*(*1*) *subt-at-subst subterm-poss-conv*)
from *p2′ Prule*(*1*,*4*) *p2* ∗ *i l as-i* **have** *get-label* ((*labeled-source* (*As*!*i*))|-*?p′*) = *Some* (α, *n*)
by (*metis fun-well-arg nth-mem*)
with *lab′* ∗ **show** *?thesis*
by (*metis p1-above-p pos-diff-def prefix-pos-diff*)
qed
qed *simp*

If a function symbol is labeled with (α, *n*), then the function symbol *n* positions above it is labeled with (α, *0*).

lemma *obtain-label-root*:
assumes *p* ∈ *poss* (*labeled-source A*)
and *get-label* ((*labeled-source A*)|-*p*) = *Some* (α, *n*)
and *A* ∈ *wf-pterm R*
shows *get-label* ((*labeled-source A*)|-(*take* (*length p* − *n*) *p*)) = *Some* (α, *0*) ∧ *take* (*length p* − *n*) *p* ∈ *poss* (*labeled-source A*)
proof−
from *assms* **have** *n*:*n* ≤ *length p*
using *label-term-max-value* **by** *blast*
with *assms* **show** *?thesis*
by (*metis* (*no-types*, *lifting*) *add.right-neutral append-take-drop-id diff-diff-cancel label-decrease length-drop poss-append-poss*)
qed

lemma *label-ctxt-apply-term*:
assumes *get-label* (*labeled-source A* |- *p*) = *l q* ∈ *poss s*
shows *get-label* (*labeled-source* ((*ctxt-of-pos-term q* (*to-pterm s*)) ⟨*A*⟩) |- (*q*@*p*)) = *l*
using *assms*(*2*) **proof**(*induct s arbitrary*:*q*)
case (*Var x*)

then have *q*:*q* = [] **by** *simp*
from *assms*(*1*) **show** *?case* **unfolding** *q* **by** *simp*
next
case (*Fun f ts*)
then show *?case* **proof**(*cases q*)
case *Nil*
from *assms*(*1*) **show** *?thesis* **unfolding** *Nil* **by** *simp*
next
case (*Cons i q′*)
with *Fun*(*2*) **have** *i*:*i* < *length ts* **and** *q′*:*q′* ∈ *poss* (*ts*!*i*) **by** *auto*
with *Fun*(*1*) **have** *get-label* (*labeled-source* (*ctxt-of-pos-term q′* (*to-pterm* (*ts*!*i*)))⟨*A*⟩ |- (*q′* @ *p*)) = *l* **by** *simp*
then show *?thesis*
unfolding *to-pterm.simps Cons ctxt-of-pos-term.simps labeled-source.simps append-Cons intp-actxt.simps subt-at.simps*
by (*metis* (*no-types*, *lifting*) *Cons-nth-drop-Suc append-take-drop-id i length-append length-map less-imp-le-nat list.size*(*4*) *nth-append-take nth-map*)
qed
qed

lemma *single-redex-at-p-label*:
assumes *p* ∈ *poss s* **and** *is-Fun* (*lhs α*)
shows *get-label* (*labeled-source* (*ll-single-redex s p α*) |-*p*) = *Some*(*α*, *0*)
proof−
from *assms*(*2*) **obtain** *f ts* **where** *f*:*lhs α* = *Fun f ts*
by *blast*
have *get-label* (*labeled-source* (*Prule α* (*map* (*to-pterm* ∘ (*λpi. s* |- (*p* @ *pi*))) (*var-poss-list* (*lhs α*))))) = *Some* (*α*, *0*)
unfolding *f labeled-source.simps label-term.simps eval-term.simps get-label.simps* **by** *simp*
then show *?thesis*
unfolding *ll-single-redex-def* **using** *label-ctxt-apply-term*[**where** *p*=[]] *assms*(*1*)
by (*smt* (*verit*) *self-append-conv subt-at.simps*(*1*))
qed

Whenever a function symbol at position *p* has label (*α*, *0*) or no label in *labeled-source A*, then we know that there exists a position *q* in *A* such that *A* |- *q* = *α As* for appropriate *As*. Moreover, taking the source of the context at position *q* must be the same as first computing the source of *A* and then taking the context at *p*.

context *left-lin*
begin
lemma *poss-labeled-source*:
assumes *p* ∈ *poss* (*labeled-source A*)
and *get-label* ((*labeled-source A*)|-*p*) = *Some* (*α*, *0*)
and *A* ∈ *wf-pterm R*
shows ∃ *q* ∈ *poss A*. *ctxt-of-pos-term p* (*source A*) = *source-ctxt* (*ctxt-of-pos-term q A*) ∧
A|-*q* = *Prule α* (*map* (*λi. A*|-(*q*@[*i*])) [*0*..<*length* (*var-rule α*)])

using *assms* **proof**(*induct A arbitrary*:*p*)
  case (*Var x*)
  then have $p = []$ **by** *simp*
  with *Var*(*2*) **have** *False* **unfolding** *labeled-source.simps* **by** *simp*
  then show *?case* **by** *blast*
next
  case (*Pfun f As*)
  then show *?case* **proof**(*cases p*)
    case (*Cons i p'*)
    with *Pfun*(*2*) **have** *ip'*:$p = i\#p'$ **and** *i*:$i < length\ As$
      by *auto*
    with *Pfun*(*2*) **have** *p'*:$p' \in poss\ (labeled\text{-}source\ (As!i))$
      by *simp*
    from *Pfun*(*3*) *i ip'* **have** $get\text{-}label\ (labeled\text{-}source\ (As!i)\ |\text{-}\ p') = Some\ (\alpha,\ 0)$
      by *simp*
    with *Pfun*(*1*,*4*) *p' i* **obtain** *q* **where** *q*:$q \in poss\ (As!i)$ **and** $ctxt\text{-}of\text{-}pos\text{-}term\ p'\ (source\ (As!i)) = source\text{-}ctxt\ (ctxt\text{-}of\text{-}pos\text{-}term\ q\ (As!i))$
      and *prule*:$(As!i)|\text{-}q = Prule\ \alpha\ (map\ (\lambda j.\ (As!i)|\text{-}(q@[j]))\ [0..<length\ (var\text{-}rule\ \alpha)])$
      using *nth-mem* **by** *blast*
    then have $ctxt\text{-}of\text{-}pos\text{-}term\ p\ (source\ (Pfun\ f\ As)) = source\text{-}ctxt\ (ctxt\text{-}of\text{-}pos\text{-}term\ (i\#q)\ (Pfun\ f\ As))$
      unfolding *ip'* **using** *i* **by**(*simp add*: *take-map drop-map*)
    then show *?thesis*
      using *q*(*1*) *i prule* **by** *fastforce*
  qed *simp*
next
  case (*Prule β As*)
  have *l*:$length\ As = length\ (var\text{-}rule\ \beta)$
    using *Prule*(*4*) **using** *wf-pterm.simps* **by** *fastforce*
  from *Prule*(*2*) **consider** $p \in fun\text{-}poss\ (labeled\text{-}lhs\ \beta)\ |\ (\exists p1\ p2\ x.\ p = p1@p2$
    $\wedge\ p1 \in poss\ (labeled\text{-}lhs\ \beta) \wedge (labeled\text{-}lhs\ \beta)|\text{-}p1 = Var\ x$
    $\wedge\ p2 \in poss\ ((\langle map\ labeled\text{-}source\ As\rangle_\beta)\ x)$
    $\wedge\ (labeled\text{-}source\ (Prule\ \beta\ As))|\text{-}p = ((\langle map\ labeled\text{-}source\ As\rangle_\beta)\ x)|\text{-}p2)$
    unfolding *labeled-source.simps* **by** (*meson poss-is-Fun-fun-poss poss-subst-choice*)
  then show *?case* **proof**(*cases*)
    case *1*
    then have $p \in fun\text{-}poss\ (lhs\ \beta)$
      by (*simp add*: *fun-poss-label-term*)
    then have $get\text{-}label\ ((labeled\text{-}source\ (Prule\ \beta\ As))|\text{-}p) = Some\ (\beta,\ length\ p)$
      unfolding *labeled-source.simps* **by** (*simp add*: *label-term-increase*)
    with *Prule*(*3*) **have** *p*:$p = []$ **and** *α*:$\alpha = \beta$ **by** *simp+*
    have $As = (map\ (\lambda i.\ Prule\ \beta\ As\ |\text{-}\ ([i]))\ [0..<length\ As])$
      by (*simp add*: *map-nth*)
    then have $As = (map\ (\lambda i.\ Prule\ \beta\ As\ |\text{-}\ ([]\ @\ [i]))\ [0..<length\ (var\text{-}rule\ \alpha)])$
      unfolding *α* **using** *l* **by** *force*
    then show *?thesis* **unfolding** *p α* **by** *auto*

next
case *2*
then obtain *p1 p2 x* **where** *p1p2*:*p* = *p1* @ *p2* **and** *x*:*p1* ∈ *poss* (*labeled-lhs* β) ∧ *labeled-lhs* β |- *p1* = *Var x*
and *p2*:*p2* ∈ *poss* ((⟨*map labeled-source As*⟩$_\beta$) *x*)
and *lab*:*labeled-source* (*Prule* β *As*) |- *p* = (⟨*map labeled-source As*⟩$_\beta$) *x* |- *p2*
by *blast*
from *Prule*(*4*) **have** *l*:*length As* = *length* (*var-rule* β)
using *wf-pterm.simps* **by** *fastforce*
from *Prule*(*4*) **have** *to-rule* β ∈ *R*
using *wf-pterm.cases* **by** *force*
with *left-lin* **have** *lin*:*linear-term* (*lhs* β)
using *left-linear-trs-def* **by** *fastforce*
from *x* **have** *p1*:*p1* ∈ *poss* (*lhs* β) **by** *simp*
then have *p1'*:*p1* ∈ *poss* ((*lhs* β) · ⟨*map source As*⟩$_\beta$) **by** *simp*
from *p1 x* **have** *x'*:*lhs* β |- *p1* = *Var x*
by (*metis label-term-to-term labeled-source-pos term-lab-to-term.simps*(*1*))
from *p1 x'* **obtain** *i* **where** *i*:*i* < *length* (*vars-term-list* (*lhs* β)) *var-poss-list* (*lhs* β) ! *i* = *p1 vars-term-list* (*lhs* β) ! *i* = *x*
by (*metis in-set-idx length-var-poss-list term.inject*(*1*) *var-poss-iff var-poss-list-sound vars-term-list-var-poss-list*)
with *lin* **have** *i'*:*i* < *length* (*var-rule* β) ∧ (*var-rule* β)!*i* = *x*
by (*metis linear-term-var-vars-term-list*)
with *l* **have** ∗:(⟨*map labeled-source As*⟩$_\beta$) *x* = *labeled-source* (*As*!*i*)
by (*metis* (*no-types*, *lifting*) *length-map lhs-subst-var-i nth-map*)
with *Prule*(*3*) *lab* **have** *get-label* ((*labeled-source* (*As*!*i*))|-*p2*) = *Some* (α, *0*)
by *simp*
with *Prule*(*1*,*4*) *p2* ∗ *i' l* **obtain** *q* **where** *q*:*q*∈*poss* (*As*!*i*) *ctxt-of-pos-term p2* (*source* (*As*!*i*)) = *source-ctxt* (*ctxt-of-pos-term q* (*As*!*i*))
(*As*!*i*) |- *q* = *Prule* α (*map* (λ*j*. (*As*!*i*) |- (*q* @ [*j*])) [*0*..<*length* (*var-rule* α)])
by (*smt* (*verit*, *ccfv-SIG*) *fun-well-arg map-eq-conv nth-mem*)
have *p1''*:(*var-poss-list* (*lhs* β) ! *length* (*take i As*)) = *p1*
using *i l* **by** (*metis id-take-nth-drop length-take length-var-poss-list lin linear-term-var-vars-term-list nth-append-length*)
have *x-sub*:*Var x* · ⟨*map source As*⟩$_\beta$ = *source* (*As*!*i*)
by (*metis* (*no-types*, *lifting*) *i' l length-map lhs-subst-var-i nth-map eval-term.simps*(*1*))
have *ctxt-of-pos-term p* (*source* (*Prule* β *As*)) = *source-ctxt* (*ctxt-of-pos-term* (*i*#*q*) (*Prule* β *As*)) **proof**−
{**fix** *y* **assume** *y* ∈ *vars-term* (*lhs* β) *y* ≠ *vars-term-list* (*lhs* β) ! *i*
then obtain *j* **where** *j*:*j* < *length* (*var-rule* β) *y* = (*var-rule* β) ! *j j* ≠ *i*
by (*metis in-set-conv-nth lin linear-term-var-vars-term-list set-vars-term-list*)
have *x*:(*vars-term-list* (*lhs* β) ! *length* (*take i As*)) = *x*
by (*metis i' id-take-nth-drop l length-take lin linear-term-var-vars-term-list nth-append-length*)
from *j* **have** (⟨*map source* (*take i As* @ *Var x* # *drop* (*Suc i*) *As*)⟩$_\beta$) *y* = *source* (*As*!*j*)
using *apply-lhs-subst-var-rule l*
by (*smt* (*verit*, *best*) *Cons-nth-drop-Suc append-Cons-nth-not-middle append-take-drop-id i' length-append length-map length-nth-simps*(*2*) *lin linear-term-var-vars-term-list*

nth-map x)
then have (⟨*map source* (*take i As* @ *Var* (*vars-term-list* (*lhs* β) ! *length* (*take i As*)) # *drop* (*Suc i*) *As*)⟩$_\beta$) *y* = (⟨*map source As*⟩$_\beta$) *y*
unfolding *x* **by** (*metis* (*no-types*, *lifting*) *j*(*1*,*2*) *l length-map lhs-subst-var-i nth-map*)
}
then have ∗:*ctxt-of-pos-term p1* (*lhs* β) $\cdot_c$ ⟨*map source As*⟩$_\beta$ =
ctxt-of-pos-term p1 (*lhs* β · ⟨*map source* (*take i As* @ *Var* (*vars-term-list* (*lhs* β) ! *length* (*take i As*)) # *drop* (*Suc i*) *As*)⟩$_\beta$)
using *i*
unfolding *ctxt-of-pos-term-subst*[*OF p1*, *symmetric*]
apply (*intro ctxt-of-pos-term-hole-subst*[*OF lin*, *of i*])
subgoal by (*metis length-var-poss-list*)
by *auto*
then show *?thesis*
unfolding *source.simps p1p2 ctxt-of-pos-term-append*[*OF p1'*] *ctxt-of-pos-term-subst*[*OF p1*] *subt-at-subst*[*OF p1*] *x' ctxt-of-pos-term.simps source-ctxt.simps Let-def x-sub q*(*2*) ∗ *p1''*
by *simp*
qed
moreover from *q*(*3*) **have** *Prule* β *As* |- (*i*#*q*) = *Prule* α (*map* (λ*j*. *Prule* β *As* |- ((*i*#*q*) @ [*j*])) [*0*..<*length* (*var-rule* α)])
by *simp*
ultimately show *?thesis*
using *i' q*(*1*) *l* **by** (*metis poss-Cons-poss term.sel*(*4*))
qed
qed

lemma *poss-labeled-source-None*:
assumes *p* ∈ *poss* (*labeled-source A*)
and *get-label* ((*labeled-source A*)|-*p*) = *None*
and *A* ∈ *wf-pterm R*
shows ∃ *q* ∈ *poss A*. *ctxt-of-pos-term p* (*source A*) = *source-ctxt* (*ctxt-of-pos-term q A*)
using *assms* **proof**(*induct A arbitrary*:*p*)
case (*Pfun f As*)
then show *?case* **proof**(*cases p*)
case (*Cons i p'*)
with *Pfun*(*2*) **have** *ip'*:*p* = *i*#*p'* **and** *i*:*i* < *length As*
by *auto*
with *Pfun*(*2*) **have** *p'*:*p'* ∈ *poss* (*labeled-source* (*As*!*i*))
by *simp*
from *Pfun*(*3*) **have** *get-label* (*labeled-source* (*As* ! *i*) |- *p'*) = *None*
unfolding *ip' labeled-source.simps* **using** *i* **by** *simp*
with *Pfun*(*1*,*4*) *p' i* **obtain** *q* **where** *q*:*q* ∈ *poss* (*As*!*i*) **and** *ctxt-of-pos-term p'* (*source* (*As*!*i*)) = *source-ctxt* (*ctxt-of-pos-term q* (*As*!*i*))
using *nth-mem* **by** *blast*
then have *ctxt-of-pos-term p* (*source* (*Pfun f As*)) = *source-ctxt* (*ctxt-of-pos-term* (*i*#*q*) (*Pfun f As*))

unfolding *ip′* **using** *i* **by**(*simp add*: *take-map drop-map*)
then show *?thesis*
using *q*(*1*) *i* **by** *fastforce*
qed *simp*
next
case (*Prule* β *As*)
have *l*:*length As* = *length* (*var-rule* β)
using *Prule*(*4*) **using** *wf-pterm.simps* **by** *fastforce*
from *Prule*(*2*) **consider** $p \in$ *fun-poss* (*labeled-lhs* β) | ($\exists$ *p1 p2 x*. *p* = *p1*@*p2*
$\wedge$ *p1* $\in$ *poss* (*labeled-lhs* β) $\wedge$ (*labeled-lhs* β)|-*p1* =
Var x
$\wedge$ *p2* $\in$ *poss* (($\langle$*map labeled-source As*$\rangle_\beta$) *x*)
$\wedge$ (*labeled-source* (*Prule* β *As*))|-*p* = (($\langle$*map*
labeled-source As$\rangle_\beta$) *x*)|-*p2*)
unfolding *labeled-source.simps* **by** (*meson poss-is-Fun-fun-poss poss-subst-choice*)
then show *?case* **proof**(*cases*)
case *1*
then have $p \in$ *fun-poss* (*lhs* β)
by (*simp add*: *fun-poss-label-term*)
then have *get-label* ((*labeled-source* (*Prule* β *As*))|-*p*) = *Some* (β, *length p*)
unfolding *labeled-source.simps* **by** (*simp add*: *label-term-increase*)
then show *?thesis*
using *Prule*(*3*) **by** *simp*
next
case *2*
then obtain *p1 p2 x* **where** *p1p2*:*p* = *p1* @ *p2* **and** *x*:*p1* $\in$ *poss* (*labeled-lhs*
β) $\wedge$ *labeled-lhs* β |- *p1* = *Var x*
and *p2*:*p2* $\in$ *poss* (($\langle$*map labeled-source As*$\rangle_\beta$) *x*)
and *lab*:*labeled-source* (*Prule* β *As*) |- *p* = ($\langle$*map labeled-source As*$\rangle_\beta$) *x* |- *p2*
by *blast*
from *Prule*(*4*) **have** *l*:*length As* = *length* (*var-rule* β)
using *wf-pterm.simps* **by** *fastforce*
from *Prule*(*4*) **have** *to-rule* $\beta \in R$
using *wf-pterm.cases* **by** *force*
with *left-lin* **have** *lin*:*linear-term* (*lhs* β)
using *left-linear-trs-def* **by** *fastforce*
from *x* **have** *p1*:*p1* $\in$ *poss* (*lhs* β) **by** *simp*
then have *p1′*:*p1* $\in$ *poss* ((*lhs* β) $\cdot$ $\langle$*map source As*$\rangle_\beta$) **by** *simp*
from *p1 x* **have** *x′*:*lhs* β |- *p1* = *Var x*
by (*metis label-term-to-term labeled-source-pos term-lab-to-term.simps*(*1*))
from *p1 x′* **obtain** *i* **where** *i*:*i* < *length* (*vars-term-list* (*lhs* β)) *var-poss-list*
(*lhs* β) ! *i* = *p1 vars-term-list* (*lhs* β) ! *i* = *x*
by (*metis in-set-idx length-var-poss-list term.inject*(*1*) *var-poss-iff var-poss-list-sound*
vars-term-list-var-poss-list)
with *lin* **have** *i′*:*i* < *length* (*var-rule* β) $\wedge$ (*var-rule* β)!*i* = *x*
by (*metis linear-term-var-vars-term-list*)
with *l* **have** $*$:($\langle$*map labeled-source As*$\rangle_\beta$) *x* = *labeled-source* (*As*!*i*)
by (*metis* (*no-types*, *lifting*) *length-map lhs-subst-var-i nth-map*)
with *Prule*(*3*) *lab* **have** *get-label* ((*labeled-source* (*As*!*i*))|-*p2*) = *None*

by *simp*
with *Prule*(*1*,*4*) *p2* ∗ *i'* *l* **obtain** *q* **where** *q*:*q*∈*poss* (*As*!*i*) *ctxt-of-pos-term p2* (*source* (*As*!*i*)) = *source-ctxt* (*ctxt-of-pos-term q* (*As*!*i*))
by (*smt* (*verit*, *ccfv-SIG*) *fun-well-arg map-eq-conv nth-mem*)
have *p1''*:(*var-poss-list* (*lhs* β) ! *length* (*take i As*)) = *p1*
using *i l* **by** (*metis id-take-nth-drop length-take length-var-poss-list lin linear-term-var-vars-term-list nth-append-length*)
have *x-sub*: *Var x* · ⟨*map source As*⟩$_\beta$ = *source* (*As*!*i*)
by (*metis* (*no-types*, *lifting*) *i' l length-map lhs-subst-var-i nth-map eval-term.simps*(*1*))
have *ctxt-of-pos-term p* (*source* (*Prule* β *As*)) = *source-ctxt* (*ctxt-of-pos-term* (*i*#*q*) (*Prule* β *As*)) **proof**−
{**fix** *y* **assume** *y* ∈ *vars-term* (*lhs* β) *y* ≠ *vars-term-list* (*lhs* β) ! *i*
then obtain *j* **where** *j*:*j* < *length* (*var-rule* β) *y* = (*var-rule* β) ! *j j* ≠ *i*
by (*metis in-set-conv-nth lin linear-term-var-vars-term-list set-vars-term-list*)
have *x*:(*vars-term-list* (*lhs* β) ! *length* (*take i As*)) = *x*
by (*metis i' id-take-nth-drop l length-take lin linear-term-var-vars-term-list nth-append-length*)
from *j* **have** (⟨*map source* (*take i As* @ *Var x* # *drop* (*Suc i*) *As*)⟩$_\beta$) *y* = *source* (*As*!*j*)
using *apply-lhs-subst-var-rule l*
by (*smt* (*verit*, *best*) *Cons-nth-drop-Suc append-Cons-nth-not-middle append-take-drop-id i' length-append length-map length-nth-simps*(*2*) *lin linear-term-var-vars-term-list nth-map x*)
then have (⟨*map source* (*take i As* @ *Var* (*vars-term-list* (*lhs* β) ! *length* (*take i As*)) # *drop* (*Suc i*) *As*)⟩$_\beta$) *y* = (⟨*map source As*⟩$_\beta$) *y*
unfolding *x* **by** (*metis* (*no-types*, *lifting*) *j*(*1*,*2*) *l length-map lhs-subst-var-i nth-map*)
}
then have ∗:*ctxt-of-pos-term p1* (*lhs* β) $\cdot_c$ ⟨*map source As*⟩$_\beta$ =
ctxt-of-pos-term p1 (*lhs* β · ⟨*map source* (*take i As* @ *Var* (*vars-term-list* (*lhs* β) ! *length* (*take i As*)) # *drop* (*Suc i*) *As*)⟩$_\beta$)
using *i*
unfolding *ctxt-of-pos-term-subst*[*OF p1*, *symmetric*]
apply (*intro ctxt-of-pos-term-hole-subst*[*OF lin*, *of i*])
subgoal by (*metis length-var-poss-list*)
by *auto*
then show *?thesis*
unfolding *source.simps p1p2 ctxt-of-pos-term-append*[*OF p1'*] *ctxt-of-pos-term-subst*[*OF p1*] *subt-at-subst*[*OF p1*] *x' ctxt-of-pos-term.simps source-ctxt.simps Let-def x-sub q*(*2*) ∗ *p1''*
by *simp*
qed
then show *?thesis*
using *i' q*(*1*) *l* **by** (*metis poss-Cons-poss term.sel*(*4*))
qed
qed *simp*
end

If we know that some part of a term does not contain labels (i.e., is coming from a simple proof term t) then we know that the label originates below some variable position of t.

lemma *labeled-source-to-pterm-subst*:
assumes *p-pos*:$p \in$ *possL* (*to-pterm* $t \cdot \sigma$) **and** *well*:$\forall x \in$ *vars-term* t. $\sigma\ x \in$ *wf-pterm R*
shows $\exists$ *p1 p2 x* γ. *p1* $\in$ *poss t* $\wedge$ $t|$-*p1* = *Var x* $\wedge$ *p1*@*p2* $\leq_p$ *p*
$\wedge$ *p2* $\in$ *possL* (σ *x*) $\wedge$ *get-label* ((*labeled-source* (σ *x*))|-*p2*) = *Some* (γ, *0*)
proof−
{**assume** *p*:*p* $\in$ *fun-poss* (*labeled-source* (*to-pterm t*))
then have *get-label* ((*labeled-source* (*to-pterm t*))|-*p*) = *None*
using *labeled-source-simple-pterm* **by** (*metis empty-iff fun-poss-imp-poss get-label-imp-labelposs*)
moreover have *get-label* ((*labeled-source* ((*to-pterm t*) $\cdot$ σ))|-*p*) = *get-label* ((*labeled-source* (*to-pterm t*))|-*p*)
by (*metis get-label-at-fun-poss-subst labeled-source-apply-subst p to-pterm-wf-pterm*)
ultimately have *False* **using** *p-pos possL-obtain-label* **by** *fastforce*
}
with *p-pos* **obtain** *p1 r x* **where** *p*:*p* = *p1*@*r* **and** *p1*:*p1* $\in$ *poss t* **and** *t*:(*labeled-source* (*to-pterm t*))|-*p1* = *Var x*
by (*smt* (*z3*) *labeled-source-apply-subst labeled-source-to-term possL-subset-poss-source poss-subst-apply-term poss-term-lab-to-term source-to-pterm subset-eq to-pterm-wf-pterm*)
then have *x*:*t*|-*p1* = *Var x*
by (*metis labeled-source-pos labeled-source-to-term source-to-pterm term-lab-to-term.simps*(*1*))
from *p-pos* **have** *r-pos*:*r* $\in$ *poss* (*labeled-source* (σ *x*))
unfolding *p* **using** *p1 t labeled-source-apply-subst*
by (*smt* (*z3*) *comp-apply labeled-source-to-term labelposs-subs-poss less-eq-pos-def less-eq-pos-simps*(*1*) *p poss-append-poss poss-term-lab-to-term source-to-pterm subset-eq eval-term.simps*(*1*) *subt-at-subst to-pterm-wf-pterm*)
from *p-pos* **obtain** γ *n* **where** *lab*:*get-label* ((*labeled-source* (σ *x*))|-*r*) = *Some* (γ, *n*)
unfolding *p labeled-source-apply-subst*[*OF to-pterm-wf-pterm*] **using** *t p1 p*
by (*smt* (*verit*, *ccfv-SIG*) *comp-apply fun-poss-imp-poss labeled-source-to-term labelposs-obtain-label labelposs-subs-fun-poss poss-term-lab-to-term source-to-pterm subset-eq eval-term.simps*(*1*) *subt-at-subst subterm-poss-conv*)
let *?p2*=*take* (*length r* − *n*) *r*
have *?p2* $\leq_p$ *r* **by** (*metis append-take-drop-id less-eq-pos-simps*(*1*))
then have *p1*@*?p2* $\leq_p$ *p* **unfolding** *p* **by** *simp*
moreover have *get-label* ((*labeled-source* (σ *x*))|-*?p2*) = *Some* (γ, *0*) $\wedge$ *?p2* $\in$ *poss* (*labeled-source* (σ *x*))
using *obtain-label-root*[*OF r-pos lab*] *well p1 x* **by** (*metis in-mono term.set-intros*(*3*) *vars-term-subt-at*)
moreover then have *?p2* $\in$ *possL* (σ *x*) **using** *get-label-imp-labelposs* **by** *blast*
ultimately show *?thesis* **using** *p1 x* **by** *blast*
qed

lemma *labelposs-subst*:
assumes *p* $\in$ *labelposs* ($t \cdot \sigma$)
shows *p* $\in$ *labelposs t* $\vee$ ($\exists$ *p1 p2 x*. *p* = *p1*@*p2* $\wedge$ *p1* $\in$ *poss t* $\wedge$ *t*|-*p1* = *Var x* $\wedge$ *p2* $\in$ *labelposs* (σ *x*))

using *assms* **proof**(*induct t arbitrary*:*p*)
case (*Fun fl ts*)
then show *?case* **proof**(*cases p*)
case *Nil*
from *Fun*(*2*) **obtain** *f l* **where** *fl* = (*f*, *Some l*)
unfolding *eval-term.simps Nil* **by** (*metis get-label.simps*(*2*) *labelposs-obtain-label subt-at.simps*(*1*) *surjective-pairing*)
then show *?thesis*
unfolding *Nil* **by** *simp*
next
case (*Cons i p′*)
from *Fun*(*2*) **have** *i*:*i* < *length ts*
unfolding *Cons eval-term.simps* **using** *labelposs-subs-poss* **by** *fastforce*
with *Fun*(*2*) **have** *p′* ∈ *labelposs* (*ts*!*i* · *σ*)
unfolding *Cons eval-term.simps* **by** (*metis* (*no-types*, *lifting*) *get-label-imp-labelposs labelposs-obtain-label labelposs-subs-poss nth-map option.simps*(*3*) *poss-Cons-poss subset-eq subt-at.simps*(*2*) *term.sel*(*4*))
with *Fun*(*1*) *i* **consider** *p′* ∈ *labelposs* (*ts*!*i*) | (∃ *p1 p2 x*. *p′* = *p1* @ *p2* ∧ *p1* ∈ *poss* (*ts*!*i*) ∧ (*ts*!*i*) |- *p1* = *Var x* ∧ *p2* ∈ *labelposs* (*σ x*))
by (*meson nth-mem*)
then show *?thesis* **proof**(*cases*)
case *1*
with *i* **show** *?thesis* **unfolding** *Cons*
by (*metis* (*no-types*, *lifting*) *get-label-imp-labelposs labelposs-obtain-label labelposs-subs-poss option.simps*(*3*) *poss-Cons-poss subsetD subt-at.simps*(*2*) *term.sel*(*4*))
next
case *2*
then obtain *p1 p2 x* **where** *p′*:*p′* = *p1* @ *p2* **and** *p1*:*p1* ∈ *poss* (*ts* ! *i*) *ts* ! *i* |- *p1* = *Var x* **and** *p2* ∈ *labelposs* (*σ x*)
by *blast*
with *i* **show** *?thesis* **unfolding** *Cons*
by (*metis append-Cons poss-Cons-poss subt-at.simps*(*2*) *term.sel*(*4*))
qed
qed
qed *simp*

lemma *set-labelposs-subst*:
labelposs (*t* · *σ*) = *labelposs t* ∪ (⋃ *i*< *length* (*vars-term-list t*). {(*var-poss-list t*!*i*)@*q* | *q*. *q* ∈ *labelposs* (*σ* (*vars-term-list t* ! *i*))}) (**is** *?ps* = *?qs*)
proof
{**fix** *p* **assume** *p* ∈ *?ps*
then consider *p* ∈ *labelposs t* | (∃ *p1 p2 x*. *p* = *p1*@*p2* ∧ *p1* ∈ *poss t* ∧ *t*|-*p1* = *Var x* ∧ *p2* ∈ *labelposs* (*σ x*))
using *labelposs-subst* **by** *blast*
then have *p* ∈ *?qs* **proof**(*cases*)
case *2*
then obtain *p1 p2 x* **where** *p* = *p1*@*p2* ∧ *p1* ∈ *poss t* ∧ *t*|-*p1* = *Var x* ∧ *p2* ∈ *labelposs* (*σ x*)
by *blast*

moreover then obtain *i* **where** *i*:$i <$ *length* (*vars-term-list t*) *vars-term-list t* ! $i = x$ *var-poss-list t* ! $i = p1$
by (*metis in-set-idx length-var-poss-list term.inject*(*1*) *var-poss-iff var-poss-list-sound vars-term-list-var-poss-list*)
ultimately have $p \in \{$*var-poss-list t* ! *i* @ *q* |*q*. $q \in$ *labelposs* (σ (*vars-term-list t* ! *i*))$\}$
by *blast*
with *i*(*1*) **show** *?thesis*
by *blast*
qed *simp*
}
then show *?ps* $\subseteq$ *?qs*
by *blast*
{fix *q* **assume** $q \in$ *?qs*
then consider $q \in$ *labelposs t* | $q \in (\bigcup i<$ *length* (*vars-term-list t*). $\{$(*var-poss-list t*!*i*)@*q* | *q*. $q \in$ *labelposs* (σ (*vars-term-list t* ! *i*))$\}$)
by *blast*
then have $q \in$ *?ps* **proof**(*cases*)
case *1*
then show *?thesis* **proof**(*induct t arbitrary*:*q*)
case (*Fun f ts*)
then show *?case* **proof**(*cases q*)
case *Nil*
with *Fun*(*2*) **obtain** *f′ lab* **where** *f′*:*f* = (*f′*, *Some lab*)
by (*metis get-label.simps*(*2*) *labelposs-obtain-label prod.exhaust-sel subt-at.simps*(*1*))
show *?thesis* **unfolding** *Nil f′* **by** *simp*
next
case (*Cons i q′*)
obtain *f′ lab* **where** *f′*:*f* = (*f′*, *lab*)
by *fastforce*
show *?thesis* **proof**(*cases lab*)
case *None*
from *Fun*(*2*) **have** *i*:$i <$ *length ts*
unfolding *f′ Cons None labelposs.simps* **by** *simp*
from *Fun*(*2*) **have** $q' \in$ *labelposs* (*ts*!*i*)
unfolding *f′ Cons None* **by** *simp*
with *Fun*(*1*) *i* **have** $q' \in$ *labelposs* (*ts*!*i* $\cdot$ σ)
by *simp*
with *i* **show** *?thesis*
unfolding *f′ Cons None eval-term.simps labelposs.simps* **by** *simp*
next
case (*Some a*)
from *Fun*(*2*) **have** *i*:$i <$ *length ts*
unfolding *f′ Cons Some labelposs.simps* **by** *simp*
from *Fun*(*2*) **have** $q' \in$ *labelposs* (*ts*!*i*)
unfolding *f′ Cons Some* **by** *simp*
with *Fun*(*1*) *i* **have** $q' \in$ *labelposs* (*ts*!*i* $\cdot$ σ)
by *simp*

with *i* **show** *?thesis*
unfolding f' *Cons Some eval-term.simps labelposs.simps* **by** *simp*
qed
qed
qed *simp*
next
case *2*
then show *?thesis* **proof**(*induct t arbitrary*:*q*)
case (*Var x*)
have *var-poss-list* (*Var x*) = [[]]
unfolding *poss-list.simps var-poss.simps* **by** *simp*
with *Var* **show** *?case* **unfolding** *vars-term-list.simps*
by (*smt* (*verit*, *ccfv-SIG*) *One-nat-def UN-iff bot-nat-0.not-eq-extremum length-0-conv length-nth-simps*(*2*) *lessThan-iff mem-Collect-eq not-less-eq nth-Cons-0 self-append-conv2 eval-term.simps*(*1*))
next
case (*Fun fl ts*)
from *Fun*(*2*) **obtain** i q' **where** q:q = *var-poss-list* (*Fun fl ts*) ! i @ q' $q' \in$ *labelposs* (σ (*vars-term-list* (*Fun fl ts*) ! i)) **and** i:i < *length* (*vars-term-list* (*Fun fl ts*))
by *blast*
then have i':i < *length* (*var-poss-list* (*Fun fl ts*))
by (*metis length-var-poss-list*)
then obtain *j r* **where** j:j < *length ts var-poss-list* (*Fun fl ts*) ! i = $j\#r$
unfolding *var-poss-list.simps* **by** (*smt* (*z3*) *add.left-neutral diff-zero length-map length-upt length-zip map-nth-eq-conv min.idem nth-concat-split nth-upt nth-zip prod.simps*(*2*))
with *i* **obtain** *x* **where** *x*:*Fun fl ts* |-($j\#r$) = *Var x*
by (*metis vars-term-list-var-poss-list*)
from *j* i' **have** $j\#r \in$ *var-poss* (*Fun fl ts*)
by (*metis nth-mem var-poss-list-sound*)
then have $r \in$ *var-poss* (*ts*!*j*)
by *simp*
then obtain i' **where** r:i'<*length* (*var-poss-list* (*ts*!*j*)) r = *var-poss-list* (*ts*!*j*) ! i'
by (*metis in-set-conv-nth var-poss-list-sound*)
moreover then have (*vars-term-list* (*Fun fl ts*) ! i) = (*vars-term-list* (*ts*!*j*) ! i')
using *x* **by** (*metis i j*(*2*) *length-var-poss-list subt-at.simps*(*2*) *term.inject*(*1*) *vars-term-list-var-poss-list*)
ultimately have $r@q' \in$ ($\bigcup$ *i*<*length* (*vars-term-list* (*ts*!*j*)). {*var-poss-list* (*ts*!*j*) ! *i* @ *q* |*q*. $q \in$ *labelposs* (σ (*vars-term-list* (*ts*!*j*) ! *i*))})
using *q*(*2*) **unfolding** *length-var-poss-list* **by** *auto*
with *Fun*(*1*) *j*(*1*) **have** *r-pos*:$r@q' \in$ *labelposs* ((*ts*!*j*) $\cdot$ σ)
using *nth-mem* **by** *blast*
obtain *f lab* **where** *f*:*fl* = (*f*, *lab*)
using *surjective-pairing* **by** *blast*
then show *?case* **proof**(*cases lab*)
case *None*

from *r-pos* **have** $j\#r@q' \in$ *labelposs* (*Fun fl ts* $\cdot$ σ)
unfolding *eval-term.simps f None labelposs.simps length-map* **using** $j(1)$
by *simp*
then show *?thesis* **unfolding** *q* $j(2)$ **by** *simp*
next
case (*Some a*)
from *r-pos* **have** $j\#r@q' \in$ *labelposs* (*Fun fl ts* $\cdot$ σ)
unfolding *eval-term.simps f Some labelposs.simps length-map* **using** $j(1)$
by *simp*
then show *?thesis* **unfolding** *q* $j(2)$ **by** *simp*
qed
qed
qed
}
then show *?qs* $\subseteq$ *?ps*
by *blast*
qed

The labeled positions in a proof term *Prule* α *As* are the function positions of *lhs* α together with all labeled positions in the arguments *As*.

lemma *possl-rule*:
assumes *length As* = *length* (*var-rule* α) *linear-term* (*lhs* α)
shows *possL* (*Prule* α *As*) = *fun-poss* (*lhs* α) $\cup$ ($\bigcup$ $i<$ (*length As*). $\{$(*var-poss-list* (*lhs* α)!i)@q $|$ q. $q \in$ *possL*(*As*!i)$\}$)
proof−
from *assms*(*1*,*2*) **have** *l*:*length* (*vars-term-list* (*labeled-lhs* α)) = *length As*
by (*metis linear-term-var-vars-term-list vars-term-list-labeled-lhs*)
have *labelposs* (*labeled-lhs* α) = *fun-poss* (*lhs* α)
by (*metis fun-poss-term-lab-to-term label-poss-labeled-lhs label-term-to-term labelposs-subs-fun-poss subsetI subset-antisym*)
moreover from *assms*(*1*,*2*) **have** $i <$ *length As* $\Longrightarrow$ ($\langle$*map labeled-source As*$\rangle_\alpha$) (*vars-term-list* (*labeled-lhs* α) ! i) = *labeled-source* (*As*!i) **for** i
using *lhs-subst-var-i linear-term-var-vars-term-list* **by** (*smt* (*verit*, *best*) *length-map nth-map vars-term-list-labeled-lhs*)
ultimately show *?thesis* **using** *set-labelposs-subst*[*of labeled-lhs* α] **unfolding** *l* *var-poss-list-labeled-lhs* **by** *force*
qed

lemma *labelposs-subs-fun-poss-source*:
assumes $p \in$ *possL A*
shows $p \in$ *fun-poss* (*source A*)
proof−
have $p \in$ *fun-poss* (*labeled-source A*)
using *assms labelposs-subs-fun-poss* **by** *blast*
then show *?thesis* **using** *fun-poss-term-lab-to-term*
by *auto*
qed

The labeled source of a context (obtained from some proof term A) applied to some proof term B is the labeled source of the context applied to the labeled source of the proof term B.

context *left-lin*
begin
lemma *label-source-ctxt*:
  assumes $A \in$ *wf-pterm R*
  and *ctxt-of-pos-term p* (*source A*) = *source-ctxt* (*ctxt-of-pos-term p′ A*)
  and $p \in$ *poss* (*source A*) **and** $p' \in$ *poss A*
 shows *labeled-source* (*ctxt-of-pos-term p′ A*)⟨*B*⟩ = (*ctxt-of-pos-term p* (*labeled-source A*))⟨*labeled-source B*⟩
  using *assms* **proof**(*induct p′ arbitrary*:*p A*)
  case *Nil*
  then have *p*:*p* = []
    using *hole-pos-ctxt-of-pos-term* **by** *force*
  then show *?case* **by** *simp*
next
  case (*Cons i p′*)
  then obtain *fl As* **where** *a*:*A* = *Fun fl As* **and** *i*:*i* < *length As* **and** *p′*:$p' \in$ *poss* (*As*!*i*)
    by (*meson args-poss*)
  then show *?case* **proof**(*cases fl*)
    case (*Inl* α)
    from *Cons*(*2*) **have** *l*:*length As* = *length* (*var-rule* α)
      unfolding *a Inl* **using** *wf-pterm.cases* **by** *auto*
    have *to-rule* $\alpha \in R$
      using *Cons*(*2*) **unfolding** *a Inl* **using** *wf-pterm.cases* **by** *force*
    with *left-lin* **have** *lin*:*linear-term* (*lhs* α)
      using *left-linear-trs-def* **by** *fastforce*
    let *?p1*=*var-poss-list* (*lhs* α) ! *i*
    from *i l lin* **have** *p1*:(*lhs* α)|-*?p1* = *Var* (*var-rule* α ! *i*)
      by (*metis linear-term-var-vars-term-list vars-term-list-var-poss-list*)
    from *i l* **have** *p1-pos*:*?p1* $\in$ *poss* (*lhs* α)
     by (*metis comp-apply length-remdups-leq length-rev length-var-poss-list nth-mem order-less-le-trans var-poss-imp-poss var-poss-list-sound*)
    let *?p2*=*hole-pos* (*source-ctxt* (*ctxt-of-pos-term p′* (*As* ! *i*)))
    have *hole-pos* (*source-ctxt* (*ctxt-of-pos-term* (*i* # *p′*) *A*)) = *?p1*@*?p2*
     unfolding *a Inl source.simps ctxt-of-pos-term.simps source-ctxt.simps Let-def hole-pos-ctxt-compose* **using** *p1-pos Cons*(*5*) *a* **by** *force*
    with *Cons*(*3*) **have** *p*:*p* = *?p1*@*?p2*
      by (*metis Cons.prems*(*3*) *hole-pos-ctxt-of-pos-term*)
    have *at-p1*:(*source A*)|-*?p1* = *source* (*As*!*i*)
      unfolding *a Inl source.simps* **using** *p1*
     by (*smt* (*verit*, *best*) *Inl i l length-map lhs-subst-var-i nth-map p1-pos eval-term.simps*(*1*) *subt-at-subst*)
    with *Cons*(*4*) **have** *p2-pos*:*?p2* $\in$ *poss* (*source* (*As*!*i*))
      unfolding *p* **by** *simp*
   from *at-p1* **have** ∗:*ctxt-of-pos-term p* (*source A*) = (*ctxt-of-pos-term ?p1* (*source A*) $\circ_c$ (*ctxt-of-pos-term ?p2* (*source* (*As* ! *i*))))
      unfolding *p* **using** *ctxt-of-pos-term-append* **using** *Cons.prems*(*3*) *p* **by** *fast-*

force
from *Cons*(*3*) **have** *ctxt-of-pos-term ?p2* (*source* (*As*!*i*)) = *source-ctxt* (*ctxt-of-pos-term* *p'* (*As*!*i*))
unfolding ∗ **unfolding** *a Inl source.simps ctxt-of-pos-term.simps source-ctxt.simps Let-def* **using** *ctxt-comp-equals Cons*(*5*) *p1-pos*
by (*smt* (*verit*, *ccfv-SIG*) *a ctxt-of-pos-term.simps*(*2*) *hole-pos.simps*(*2*) *hole-pos-ctxt-of-pos-term list.inject poss-imp-subst-poss*)
with *Cons*(*1*,*2*) *i p2-pos p' a Inl* **have** *IH*:*labeled-source* (*ctxt-of-pos-term p'* (*As*!*i*))⟨*B*⟩ = (*ctxt-of-pos-term ?p2* (*labeled-source* (*As*!*i*)))⟨*labeled-source B*⟩
by (*meson fun-well-arg nth-mem*)
then have *list-IH*:*map labeled-source* (*take i As* @ (*ctxt-of-pos-term p'* (*As* ! *i*))⟨*B*⟩ # *drop* (*Suc i*) *As*) =
map labeled-source (*take i As*) @ (*ctxt-of-pos-term ?p2* (*labeled-source* (*As* ! *i*)))⟨*labeled-source B*⟩ # *map labeled-source* (*drop* (*Suc i*) *As*)
using *i* **by** *fastforce*
from *lin* **have** *lin'*:*linear-term* (*labeled-lhs* α)
using *linear-label-term* **by** *blast*
from *p1-pos* **have** *p1-pos*:*?p1* ∈ *poss* (*labeled-lhs* α)
by *simp*
from *p1* **have** *x*:*labeled-lhs* α |- *var-poss-list* (*lhs* α) ! *i* = *Var* (*var-rule* α ! *i*)
by (*metis label-term-to-term p1-pos poss-term-lab-to-term var-label-term*)
have (⟨*map labeled-source As*⟩$_\alpha$)((*var-rule* α ! *i*) := (*ctxt-of-pos-term ?p2* ((⟨*map labeled-source As*⟩$_\alpha$) (*var-rule* α ! *i*)))⟨*labeled-source B*⟩)
= ⟨(*take i* (*map labeled-source As*)) @ (*ctxt-of-pos-term ?p2* (*labeled-source* (*As*!*i*)))⟨*labeled-source B*⟩ # (*drop* (*Suc i*) (*map labeled-source As*))⟩$_\alpha$
using *i* **by** (*smt* (*verit*, *best*) *Cons.prems*(*4*) *a ctxt-of-pos-term.simps*(*2*) *hole-pos.simps*(*2*) *hole-pos-ctxt-of-pos-term id-take-nth-drop l length-map lhs-subst-upd lhs-subst-var-i list.inject nth-map take-map*)
then show *?thesis* **unfolding** *a Inl ctxt-of-pos-term.simps labeled-source.simps intp-actxt.simps p list-IH*
using *replace-at-append-subst*[*OF lin' p1-pos x*] **by** (*smt* (*verit*) *drop-map take-map*)
next
case (*Inr f*)
from *Cons*(*3*,*4*,*5*) **obtain** *p2* **where** *p*:*p* = *i*#*p2* **and** *p2*:*p2* ∈ *poss* (*source* (*As*!*i*)) **and** *ctxt*:*ctxt-of-pos-term p2* (*source* (*As*!*i*)) = *source-ctxt* (*ctxt-of-pos-term* *p'* (*As*!*i*))
unfolding *a Inr source.simps ctxt-of-pos-term.simps source-ctxt.simps* **by** (*smt* (*verit*, *best*) *Cons.prems*(*2*) *Cons.prems*(*3*) *Inr a actxt.inject ctxt-of-pos-term.simps*(*2*) *i nth-map source-poss*)
from *Cons*(*1*,*2*) *ctxt p2 p'* **have** *IH*:*labeled-source* (*ctxt-of-pos-term p'* (*As*!*i*))⟨*B*⟩ = (*ctxt-of-pos-term p2* (*labeled-source* (*As*!*i*)))⟨*labeled-source B*⟩
using *a i nth-mem* **by** *blast*
then have *list-IH*:*map labeled-source* (*take i As* @ (*ctxt-of-pos-term p'* (*As* ! *i*))⟨*B*⟩ # *drop* (*Suc i*) *As*) =
map labeled-source (*take i As*) @ (*ctxt-of-pos-term p2* (*labeled-source* (*As* ! *i*)))⟨*labeled-source B*⟩ # *map labeled-source* (*drop* (*Suc i*) *As*)
using *i* **by** *fastforce*
show *?thesis* **unfolding** *a Inr ctxt-of-pos-term.simps p labeled-source.simps*

intp-actxt.simps list-IH
by (*simp add*: *drop-map i take-map*)
qed
qed
end

lemma *labeled-ctxt-above*:
assumes $p \in$ *poss A* **and** $r \in$ *poss A* **and** $\neg\ p \leq_p r$
shows *get-label* ((*ctxt-of-pos-term p A*)$\langle$*labeled-source B*$\rangle$ |-*r*) = *get-label* (*A* |-*r*)
using *assms* **proof**(*induct A arbitrary*:*r p*)
case (*Fun f As*)
then have $p \neq []$
by *blast*
with *Fun*(*2*) **obtain** *i p′* **where** *i*:*i* < *length As* **and** *p′*:*p′* $\in$ *poss* (*As*!*i*) **and** *p*:*p* = *i*#*p′*
by *auto*
from *Fun*(*4*) **consider** $r <_p p$ | $r \perp p$
using *parallel-pos* **by** *fastforce*
then show *?case* **proof**(*cases*)
case *1*
then show *?thesis* **proof**(*cases r*)
case *Nil*
show *?thesis* **unfolding** *p Nil* **by** *simp*
next
case (*Cons j r′*)
from *1* **have** *j*:*j* = *i*
unfolding *p Cons* **by** *simp*
with *Fun*(*1*) **have** *get-label* ((*ctxt-of-pos-term p′* (*As*!*i*))$\langle$*labeled-source B*$\rangle$ |- *r′*) = *get-label* ((*As*!*i*) |- *r′*)
using *i p′ Fun*(*3*,*4*) **unfolding** *Cons j p* **by** *simp*
then show *?thesis*
unfolding *Cons p subt-at.simps ctxt-of-pos-term.simps intp-actxt.simps* **by** (*metis i j nat-less-le nth-append-take*)
qed
next
case *2*
then obtain *j r′* **where** *r*:*r* = *j*#*r′*
unfolding *p* **by** (*metis parallel-pos.elims*(*2*))
then show *?thesis* **proof**(*cases i* = *j*)
case *True*
from *Fun*(*1*) *2 i* **have** *get-label* ((*ctxt-of-pos-term p′* (*As*!*i*))$\langle$*labeled-source B*$\rangle$ |- *r′*) = *get-label* ((*As*!*i*) |- *r′*)
using *Fun.prems*(*2*) *Fun.prems*(*3*) *True p p′ r* **by** *force*
then show *?thesis* **using** *p r True*
by (*metis 2 Fun.prems*(*1*) *Fun.prems*(*2*) *parallel-pos parallel-replace-at-subt-at*)

next
case *False*
then show *?thesis*

unfolding *p r subt-at.simps ctxt-of-pos-term.simps intp-actxt.simps* **by** (*metis i nth-list-update upd-conv-take-nth-drop*)
qed
qed
qed *simp*

The labeled positions of a context (obtained from some proof term A) applied to some proof term B are the labeled positions of the context together with the labeled positions of the proof term B.

context *left-lin*
begin
lemma *label-ctxt*:
assumes $A \in$ *wf-pterm R*
and *ctxt-of-pos-term p* (*source A*) = *source-ctxt* (*ctxt-of-pos-term p′ A*)
and $p \in$ *poss* (*source A*) **and** $p' \in$ *poss A*
shows *possL* (*ctxt-of-pos-term p′ A*)⟨*B*⟩ = $\{q.\ q \in$ *possL A* $\land \neg p \leq_p q\} \cup \{p@q|$ $q.\ q \in$ *possL B*$\}$
using *assms* **proof**(*induct p′ arbitrary*:*p A*)
case *Nil*
then have *p*:*p* = []
using *hole-pos-ctxt-of-pos-term* **by** *force*
then have $\{q \in$ *possL A*. $\neg\ p \leq_p q\} = \{\}$
by *simp*
then show *?case*
unfolding *Nil ctxt-of-pos-term.simps p* **by** *simp*
next
case (*Cons i p′*)
then obtain *fl As* **where** *a*:*A* = *Fun fl As* **and** *i*:*i* < *length As* **and** *p′*:*p′* ∈ *poss* (*As*!*i*)
by (*meson args-poss*)
then show *?case* **proof**(*cases fl*)
case (*Inl α*)
from *Cons*(*2*) **have** *l*:*length As* = *length* (*var-rule α*)
unfolding *a Inl* **using** *wf-pterm.cases* **by** *auto*
have *to-rule α* ∈ *R*
using *Cons*(*2*) **unfolding** *a Inl* **using** *wf-pterm.cases* **by** *force*
with *left-lin* **have** *lin*:*linear-term* (*lhs α*)
using *left-linear-trs-def* **by** *fastforce*
let *?p1*=*var-poss-list* (*lhs α*) ! *i*
from *i l lin* **have** *p1*:(*lhs α*)|-*?p1* = *Var* (*var-rule α* ! *i*)
by (*metis linear-term-var-vars-term-list vars-term-list-var-poss-list*)
from *i l* **have** *p1-pos*:*?p1* ∈ *poss* (*lhs α*)
by (*metis comp-apply length-remdups-leq length-rev length-var-poss-list nth-mem order-less-le-trans var-poss-imp-poss var-poss-list-sound*)
let *?p2*=*hole-pos* (*source-ctxt* (*ctxt-of-pos-term p′* (*As* ! *i*)))
have *hole-pos* (*source-ctxt* (*ctxt-of-pos-term* (*i* # *p′*) *A*)) = *?p1*@*?p2*
unfolding *a Inl source.simps ctxt-of-pos-term.simps source-ctxt.simps Let-def hole-pos-ctxt-compose* **using** *p1-pos Cons*(*5*) *a* **by** *force*
with *Cons*(*3*) **have** *p*:*p* = *?p1*@*?p2*

by (*metis Cons.prems*(*3*) *hole-pos-ctxt-of-pos-term*)
have *at-p1*:(*source A*)|-*?p1* = *source* (*As*!*i*)
unfolding *a Inl source.simps* **using** *p1*
by (*smt* (*verit*, *best*) *Inl i l length-map lhs-subst-var-i nth-map p1-pos eval-term.simps*(*1*) *subt-at-subst*)
with *Cons*(*4*) **have** *p2-pos*:*?p2* ∈ *poss* (*source* (*As*!*i*))
unfolding *p* **by** *simp*
from *at-p1* **have** ∗:*ctxt-of-pos-term p* (*source A*) = (*ctxt-of-pos-term ?p1* (*source A*) $\circ_c$ (*ctxt-of-pos-term ?p2* (*source* (*As* ! *i*))))
unfolding *p* **using** *ctxt-of-pos-term-append* **using** *Cons.prems*(*3*) *p* **by** *fastforce*
from *Cons*(*3*) **have** *ctxt-of-pos-term ?p2* (*source* (*As*!*i*)) = *source-ctxt* (*ctxt-of-pos-term p′* (*As*!*i*))
unfolding ∗ **unfolding** *a Inl source.simps ctxt-of-pos-term.simps source-ctxt.simps Let-def* **using** *ctxt-comp-equals Cons*(*5*) *p1-pos*
by (*smt* (*verit*, *ccfv-SIG*) *a ctxt-of-pos-term.simps*(*2*) *hole-pos.simps*(*2*) *hole-pos-ctxt-of-pos-term list.inject poss-imp-subst-poss*)
with *Cons*(*1*,*2*) *i p2-pos p′ a Inl* **have** *IH*:*possL* (*ctxt-of-pos-term p′* (*As*!*i*))⟨*B*⟩ = {*q* ∈ *possL* (*As*!*i*). ¬ *?p2* $\leq_p$ *q*} ∪ {*?p2* @ *q* |*q*. *q* ∈ *possL B*}
by (*meson fun-well-arg nth-mem*)
let *?a1*=*fun-poss* (*lhs* α)
let *?a2*=(⋃*j* ∈ {*k*. *k* < *length As* ∧ *k* ≠ *i*}. {(*var-poss-list* (*lhs* α)!*j*)@*q* | *q*. *q* ∈ *possL*(*As*!*j*)})
let *?a3*={*?p1*@*q* | *q*. *q* ∈ *possL* (*As*!*i*) ∧ ¬ *?p2* $\leq_p$ *q*}
let *?a4*={*?p1* @ *?p2* @ *q* |*q*. *q* ∈ *possL B*}
let *?b1*={*q* ∈ *possL A*. ¬ *p* $\leq_p$ *q*}
have *?a1* ∪ *?a2* ∪ *?a3* = *?b1* **proof**
{**fix** *x* **assume** *x*:*x* ∈ *?a1*
then have ¬ *?p1* $\leq_p$ *x*
by (*metis append.right-neutral fun-poss-append-poss fun-poss-fun-conv fun-poss-imp-poss p1 prefix-pos-diff term.distinct*(*1*))
then have ¬ *p* $\leq_p$ *x*
unfolding *p* **using** *less-eq-pos-simps*(*1*) *order-pos.order.trans* **by** *blast*
with *x* **have** *x* ∈ *?b1*
unfolding *a Inl* **using** *possl-rule l lin* **by** *auto*
} **moreover** {**fix** *x* **assume** *x* ∈ *?a2*
then obtain *j q* **where** *j*:*j* < *length As j* ≠ *i* **and** *q*:*q* ∈ *possL* (*As* ! *j*) **and** *x*:*x* = *var-poss-list* (*lhs* α) ! *j* @ *q*
by *blast*
from *j* **have** *j′*:*j* < *length* (*var-poss-list* (*lhs* α))
using *l lin* **by** (*metis length-var-poss-list linear-term-var-vars-term-list*)
with *j*(*2*) **have** *?p1* ≠ (*var-poss-list* (*lhs* α)) !*j*
by (*metis* (*mono-tags*, *lifting*) *distinct-remdups distinct-rev i j*(*1*) *l lin linear-term-var-vars-term-list nth-eq-iff-index-eq o-apply term.inject*(*1*) *vars-term-list-var-poss-list*)
with *j′* **have** *?p1* ⊥ *var-poss-list* (*lhs* α) ! *j*
using *var-poss-parallel* **by** (*metis nth-mem p1 p1-pos var-poss-iff var-poss-list-sound*)
then have ¬ *p* $\leq_p$ *x*
unfolding *p x* **using** *less-eq-pos-simps*(*1*) *order-pos.order-trans pos-less-eq-append-not-parallel* **by** *blast*

then have $x \in ?b1$
unfolding a *Inl possl-rule*[*OF l lin*] x **using** $j(1)$ q **by** *blast*
} **moreover** {**fix** x **assume** $x \in ?a3$
then obtain q **where** $x{:}x = ?p1@q$ $q \in possL$ $(As ! i)$ $\neg$ $?p2 \leq_p q$
by *blast*
from $x(3)$ **have** $\neg$ $p \leq_p x$
unfolding p $x(1)$ **using** *less-eq-pos-simps*(2) **by** *blast*
with $x(2)$ **have** $x \in ?b1$
unfolding a *Inl possl-rule*[*OF l lin*] $x(1)$ **using** i **by** *auto*
}
ultimately show $?a1 \cup ?a2 \cup ?a3 \subseteq ?b1$ **by** *blast*
{**fix** x **assume** $b1{:}x \in ?b1$
then consider $x \in$ *fun-poss* $(lhs\ \alpha)$ | $x \in (\bigcup i{<}length\ As.\ \{var\text{-}poss\text{-}list\ (lhs\ \alpha)\ !\ i\ @\ q\ |q.\ q \in possL\ (As\ !\ i)\})$
unfolding a *Inl possl-rule*[*OF l lin*] **by** *blast*
then have $x \in ?a1 \cup ?a2 \cup ?a3$ **proof**(*cases*)
case *2*
then obtain j q **where** $j{:}j < length\ As$ **and** $x{:}x = var\text{-}poss\text{-}list\ (lhs\ \alpha)\ !\ j\ @\ q$ **and** $q{:}q \in possL\ (As!j)$
by *blast*
then show *?thesis* **proof**(*cases* $j = i$)
case *True*
from $b1$ **have** $\neg$ $?p2 \leq_p q$
unfolding p x *True* **using** *less-eq-pos-simps*(2) **by** *blast*
then show *?thesis* **using** j x q **by** *auto*
qed *auto*
qed *simp*
}
then show $?b1 \subseteq ?a1 \cup ?a2 \cup ?a3$ **by** *blast*
qed
moreover have $possL\ (ctxt\text{-}of\text{-}pos\text{-}term\ (i\ \#\ p')\ A)\langle B\rangle = ?a1 \cup ?a2 \cup ?a3 \cup ?a4$ **proof**−
from l i **have** $l'{:}length\ (take\ i\ As\ @\ (ctxt\text{-}of\text{-}pos\text{-}term\ p'\ (As\ !\ i))\langle B\rangle\ \#\ drop\ (Suc\ i)\ As) = length\ (var\text{-}rule\ \alpha)$
by *simp*
have $set{:}\{j.\ j < length\ As\} = \{j.\ j < length\ As \wedge j \neq i\} \cup \{i\}$
using i *Collect-disj-eq* **by** *auto*
let $?args{=}(take\ i\ As\ @\ (ctxt\text{-}of\text{-}pos\text{-}term\ p'\ (As\ !\ i))\langle B\rangle\ \#\ drop\ (Suc\ i)\ As)$
{**fix** j **assume** $j < length\ As \wedge j \neq i$
with i **have** $?args\ !\ j = As!j$
by (*meson nat-less-le nth-append-take-drop-is-nth-conv*)
} **moreover have** $?args!i = (ctxt\text{-}of\text{-}pos\text{-}term\ p'\ (As\ !\ i))\langle B\rangle$ **using** i
by (*simp add*: *nth-append-take*)
moreover from *set* **have** $(\bigcup j{<}length\ As.\ \{var\text{-}poss\text{-}list\ (lhs\ \alpha)\ !\ j\ @\ q\ |q.\ q \in possL\ (?args\ !\ j)\}) =$
$(\bigcup j \in \{j.\ j < length\ As \wedge j \neq i\}.\ \{var\text{-}poss\text{-}list\ (lhs\ \alpha)\ !\ j\ @\ q\ |q.\ q \in possL\ (?args\ !\ j)\}) \cup \{?p1\ @\ q\ |q.\ q \in possL\ (?args!i)\}$
by *force*
ultimately have $(\bigcup j{<}length\ As.\ \{var\text{-}poss\text{-}list\ (lhs\ \alpha)\ !\ j\ @\ q\ |q.\ q \in possL$

(*?args* ! *j*)}) =
($\bigcup$*j* ∈ {*j*. *j* < *length As* ∧ *j* ≠ *i*}. {*var-poss-list* (*lhs* α) ! *j* @ *q* |*q*. *q* ∈ *possL* (*As* ! *j*)}) ∪ {*?p1* @ *q* |*q*. *q* ∈ *possL* (*ctxt-of-pos-term p′* (*As* ! *i*))⟨*B*⟩}
by *simp*
moreover have *possL* (*ctxt-of-pos-term* (*i* # *p′*) *A*)⟨*B*⟩ = *fun-poss* (*lhs* α) ∪
($\bigcup$*j*<*length As*. {*var-poss-list* (*lhs* α) ! *j* @ *q* |*q*. *q* ∈ *possL* (*?args* ! *j*)})
unfolding *a Inl ctxt-of-pos-term.simps intp-actxt.simps* **using** *possl-rule*[*OF l′ lin*] *i* **by** *force*
ultimately show *?thesis* **unfolding** *IH* **by** *auto*
qed
ultimately show *?thesis* **using** *p* **by** *force*
next
case (*Inr f*)
from *Cons*(*3*,*4*,*5*) **obtain** *p2* **where** *p*:*p* = *i*#*p2* **and** *p2* ∈ *poss* (*source* (*As*!*i*)) **and** *ctxt*:*ctxt-of-pos-term p2* (*source* (*As*!*i*)) = *source-ctxt* (*ctxt-of-pos-term p′* (*As*!*i*))
unfolding *a Inr source.simps ctxt-of-pos-term.simps source-ctxt.simps* **by** (*smt* (*verit*, *best*) *Cons.prems*(*2*) *Cons.prems*(*3*) *Inr a actxt.inject ctxt-of-pos-term.simps*(*2*) *i nth-map source-poss*)
with *Cons*(*1*,*2*) *i p′* **have** *IH*:*possL* (*ctxt-of-pos-term p′* (*As*!*i*))⟨*B*⟩ = {*q* ∈ *possL* (*As*!*i*). ¬ *p2* $\leq_p$ *q*} ∪ {*p2* @ *q* |*q*. *q* ∈ *possL B*}
unfolding *a Inr* **by** (*meson fun-well-arg nth-mem*)
let *?a2*=($\bigcup$*j* ∈ {*k*. *k* < *length As* ∧ *k* ≠ *i*}. {*j* # *q* | *q*. *q* ∈ *possL*(*As*!*j*)})
let *?a3*={*i*#*q* | *q*. *q* ∈ *possL* (*As*!*i*) ∧ ¬ *p2* $\leq_p$ *q*}
let *?a4*={*i* # *p2* @ *q* |*q*. *q* ∈ *possL B*}
let *?b1*={*q* ∈ *possL A*. ¬ *p* $\leq_p$ *q*}
have *?a2* ∪ *?a3* = *?b1* **proof**
{**fix** *x* **assume** *x* ∈ *?a2*
then obtain *j q* **where** *j*:*j* < *length As j* ≠ *i* **and** *q*:*q* ∈ *possL* (*As* ! *j*) **and** *x*:*x* = *j* # *q*
by *blast*
from *j q* **have** *j*#*q* ∈ *possL A*
unfolding *a Inr* **by** *simp*
then have *x* ∈ *?b1*
unfolding *x p* **using** *j*(*2*) **by** *simp*
} **moreover** {**fix** *x* **assume** *x* ∈ *?a3*
then obtain *q* **where** *x*:*x* = *i*#*q q* ∈ *possL* (*As* ! *i*) ¬ *p2* $\leq_p$ *q*
by *blast*
from *x*(*3*) **have** ¬ *p* $\leq_p$ *x*
unfolding *p x*(*1*) **using** *less-eq-pos-simps*(*2*) **by** *simp*
with *x*(*2*) **have** *x* ∈ *?b1*
unfolding *a Inr x*(*1*) **using** *i* **by** *auto*
}
ultimately show *?a2* ∪ *?a3* ⊆ *?b1* **by** *blast*
{**fix** *x* **assume** *b1*:*x* ∈ *?b1*
then have *x* ∈ *possL A*
by *simp*
then obtain *j q* **where** *j*:*j* < *length As* **and** *x*:*x* = *j* # *q* **and** *q*:*q* ∈ *possL*

$(As!j)$
unfolding *a Inr labeled-source.simps labelposs.simps length-map* **by** *force*
then have $x \in ?a2 \cup ?a3$ **proof**(*cases* $j = i$)
case *True*
with *b1* **have** $\neg$ $p2 \leq_p q$
unfolding *p x* **using** *less-eq-pos-simps(2)* **by** *simp*
then show *?thesis* **using** *j x q b1* **by** *auto*
qed *simp*
}
then show $?b1 \subseteq ?a2 \cup ?a3$ **by** *blast*
qed
moreover have *possL* (*ctxt-of-pos-term* $(i \# p')$ $A)\langle B\rangle = ?a2 \cup ?a3 \cup ?a4$
proof−
have *l*:*length* (*take i As* @ (*ctxt-of-pos-term* p' (*As* ! *i*))$\langle B\rangle$ # *drop* (*Suc i*) *As*) = *length As*
using *i* **by** *simp*
{fix *j* **assume** $j <$ *length As*
then have (*map labeled-source* (*take i As* @ (*ctxt-of-pos-term* p' (*As* ! *i*))$\langle B\rangle$ # *drop* (*Suc i*) *As*) ! *j*) = *labeled-source* ((*take i As* @ (*ctxt-of-pos-term* p' (*As* ! *i*))$\langle B\rangle$ # *drop* (*Suc i*) *As*) ! *j*)
using *nth-map l* **by** *metis*
}note *map-lab=this*
have *set*:$\{j.\ j <$ *length As*$\} = \{j.\ j <$ *length As* $\wedge j \neq i\} \cup \{i\}$
using *i Collect-disj-eq* **by** *auto*
let *?args*=(*take i As* @ (*ctxt-of-pos-term* p' (*As* ! *i*))$\langle B\rangle$ # *drop* (*Suc i*) *As*)
{fix *j* **assume** $j <$ *length As* $\wedge j \neq i$
with *i* **have** *?args* ! j = *As*!*j*
by (*meson nat-less-le nth-append-take-drop-is-nth-conv*)
} moreover have *?args*!*i* = (*ctxt-of-pos-term* p' (*As* ! *i*))$\langle B\rangle$ **using** *i*
by (*simp add*: *nth-append-take*)
moreover from *set* **have** $(\bigcup j<$*length As*. $\{j \# q\ |q.\ q \in$ *possL* (*?args* ! *j*)$\})$ =
$(\bigcup j \in \{j.\ j <$ *length As* $\wedge j \neq i\}$. $\{j \# q\ |q.\ q \in$ *possL* (*?args* ! *j*)$\})$ $\cup$ $\{i \# q\ |q.\ q \in$ *possL* (*?args*!*i*)$\}$
by *force*
ultimately have $(\bigcup j<$*length As*. $\{j \# q\ |q.\ q \in$ *possL* (*?args* ! *j*)$\})$ =
$(\bigcup j \in \{j.\ j <$ *length As* $\wedge j \neq i\}$. $\{j \# q\ |q.\ q \in$ *possL* (*As* ! *j*)$\})$ $\cup$ $\{i \# q\ |q.\ q \in$ *possL* (*ctxt-of-pos-term* p' (*As* ! *i*))$\langle B\rangle\}$
by *simp*
moreover have *possL* (*ctxt-of-pos-term* $(i \# p')$ $A)\langle B\rangle$ = $(\bigcup j<$*length As*. $\{j \# q\ |q.\ q \in$ *possL* (*?args* ! *j*)$\})$
unfolding *a Inr ctxt-of-pos-term.simps intp-actxt.simps labeled-source.simps labelposs.simps length-map l* **using** *map-lab* **by** *force*
ultimately show *?thesis* **unfolding** *IH* **by** *auto*
qed
ultimately show *?thesis* **using** *p* **by** *force*
qed
qed

lemma *single-redex-possL*:
assumes *to-rule* $\alpha \in R$ $p \in$ *poss* s
shows *possL* (*ll-single-redex* s p α) = $\{p$ @ q $|q.$ $q \in$ *fun-poss* (*lhs* α)$\}$
proof−
let *?*Δ=*ll-single-redex* s p α
have $*$:*possL* (*Prule* α (*map* (*to-pterm* $\circ$ (λ*pi*. s|-(p@*pi*))) (*var-poss-list* (*lhs* α)))) = *labelposs* (*labeled-lhs* α)
proof−
{**fix** x
have *labelposs* (($\langle$*map labeled-source* (*map* (*to-pterm* $\circ$ (λ*pi*. s |- (p @ *pi*))) (*var-poss-list* (*lhs* α)))$\rangle_\alpha$) x) = {}
by (*smt* (*verit*) *comp-apply labeled-source-simple-pterm labelposs.simps*(*1*) *length-map lhs-subst-not-var-i lhs-subst-var-i map-nth-eq-conv*)
}
then show *?thesis* **unfolding** *labeled-source.simps labelposs-apply-subst* **by** *blast*
qed
have *possL ?*Δ = $\{q \in$ *possL* (*to-pterm* s). $\neg$ $p \leq_p q\} \cup \{p$ @ q $|q.$ $q \in$ *possL* (*Prule* α (*map* (*to-pterm* $\circ$ (λ*pi*. s|-(p@*pi*))) (*var-poss-list* (*lhs* α))))$\}$
using *label-ctxt assms* **by** (*simp add*: *ll-single-redex-def p-in-poss-to-pterm source-ctxt-to-pterm*)
also have ...= $\{p$ @ q $|q.$ $q \in$ *possL* (*Prule* α (*map* (*to-pterm* $\circ$ (λ*pi*. s|-(p@*pi*))) (*var-poss-list* (*lhs* α))))$\}$
using *labeled-source-simple-pterm* **by** *auto*
also have ...= $\{p$ @ q $|q.$ $q \in$ *labelposs* (*labeled-lhs* α)$\}$
unfolding $*$ **by** *simp*
finally show *?thesis*
using *label-poss-labeled-lhs labelposs-subs-fun-poss* **by** *fastforce*
qed

end

lemma *labeled-poss-in-lhs*:
assumes *p-pos*:$p \in$ *poss* (*source* (*Prule* α *As*)) **and** *well*:*Prule* α *As* $\in$ *wf-pterm* R
and *get-label* ((*labeled-source* (*Prule* α *As*))|-p) = *Some* (α, *length* p) *is-Fun* (*lhs* α)
shows $p \in$ *fun-poss* (*lhs* α)
proof−
from *p-pos* **consider** $p \in$ *fun-poss* (*lhs* α) | $\exists$ *p1 p2 x*. p = *p1* @ *p2* $\wedge$ *p1* $\in$ *poss* (*lhs* α) $\wedge$ (*lhs* α)|-*p1* = *Var* x $\wedge$ *p2* $\in$ *poss* (($\langle$*map source As*$\rangle_\alpha$) x)
unfolding *source.simps* **using** *poss-subst-apply-term* **by** *metis*
then show *?thesis* **proof**(*cases*)
case *2*
then obtain *p1 p2 x* **where** p:p = *p1* @ *p2* **and** *p1*:*p1* $\in$ *poss* (*lhs* α) (*lhs* α)|-*p1* = *Var* x **and** *p2*:*p2* $\in$ *poss* (($\langle$*map source As*$\rangle_\alpha$) x)
by *blast*
then obtain i **where** i:i < *length* (*var-rule* α) *var-rule* α!i = x

by (*metis in-set-conv-nth set-vars-term-list subt-at-imp-supteq subteq-Var-imp-in-vars-term vars-term-list-vars-distinct*)
from *p1* **have** *p1-pos'*:*p1* ∈ *poss* (*labeled-lhs* α)
by *simp*
from *p1* **have** *p1-pos*:*p1* ∈ *poss* (*labeled-lhs* α · ⟨*map labeled-source As*⟩$_\alpha$)
by (*metis labeled-source.simps*(*3*) *labeled-source-to-term p p-pos poss-append-poss poss-term-lab-to-term*)
from *p1* **have** *x*:*labeled-lhs* α |-*p1* = *Var x*
by (*metis fun-poss-term-lab-to-term label-term-to-term labeled-source-pos poss-simps*(*4*) *poss-term-lab-to-term term.sel*(*1*) *term-lab-to-term.simps*(*1*) *var-poss-iff*)
from *well* **have** *l*:*length As* = *length* (*var-rule* α)
using *wf-pterm.cases* **by** *auto*
with *well i* **have** *asi*:*As*!*i* ∈ *wf-pterm R*
by (*metis fun-well-arg nth-mem*)
from *l* **have** *lab*:*labeled-source* (*Prule* α *As*) |-*p* = *labeled-source* (*As*!*i*) |-*p2*
unfolding *p labeled-source.simps subt-at-append*[*OF p1-pos*] *subt-at-subst*[*OF p1-pos'*] *x* **using** *i*
by (*metis* (*no-types*, *lifting*) *length-map lhs-subst-var-i nth-map eval-term.simps*(*1*))
moreover from *assms*(*4*) *p1* **have** *length p2* < *length p*
unfolding *p* **by** *auto*
moreover from *p2* **have** *p2* ∈ *poss* (*labeled-source* (*As*!*i*))
using *l i* **by** (*metis* (*no-types*, *lifting*) *labeled-source-to-term length-map lhs-subst-var-i nth-map poss-term-lab-to-term*)
ultimately have *False* **using** *assms*(*3*) *asi* **by** (*simp add*: *label-term-max-value leD*)
then show *?thesis* **by** *simp*
qed *simp*
qed

context *left-lin-no-var-lhs*
begin
lemma *get-label-Prule*:
assumes *Prule* α *As* ∈ *wf-pterm R* **and** *p* ∈ *poss* (*source* (*Prule* α *As*)) **and** *get-label* (*labeled-source* (*Prule* α *As*) |- *p*) = *Some* (β, *0*)
shows (*p* = [] ∧ α = β) ∨
(∃ *p1 p2 i*. *p* = *p1*@*p2* ∧ *i* < *length As* ∧ *var-poss-list* (*lhs* α)!*i* = *p1* ∧
p2 ∈ *poss* (*source* (*As*!*i*)) ∧ *get-label* (*labeled-source* (*As*!*i*)|-*p2*) = *Some* (β, *0*))
proof−
from *assms*(*1*) **have** α:*to-rule* α ∈ *R*
using *wf-pterm.simps* **by** *fastforce*
with *no-var-lhs* **obtain** *f ts* **where** *lhs*:*lhs* α = *Fun f ts* **by** *fastforce*
from *assms*(*1*) **have** *l1*:*length* (*var-rule* α) = *length As*
using *wf-pterm.cases* **by** *force*
then have *l2*:*length* (*var-poss-list* (*lhs* α)) = *length As*
using *left-lin.length-var-rule*[*OF left-lin-axioms* α] **by** (*simp add*: *length-var-poss-list*)

from *left-lin* **have** *var-rule*:*var-rule* α = *vars-term-list* (*lhs* α)
using α *left-linear-trs-def linear-term-var-vars-term-list* **by** *fastforce*

then show *?thesis* **proof**(*cases p*=[])
case *True*
from *assms*(*3*) **have** $\beta = \alpha$ **unfolding** *True labeled-source.simps lhs label-term.simps eval-term.simps subt-at.simps* **by** *simp*
then show *?thesis* **unfolding** *True* **by** *simp*
next
case *False*
from *assms*(*3*) **have** *possL*:*p* ∈ *possL* (*Prule* α *As*)
by (*metis assms*(*2*) *get-label-imp-labelposs labeled-source-to-term option.distinct*(*1*) *poss-term-lab-to-term*)
{**assume** *p* ∈ *fun-poss* (*lhs* α)
then have *get-label* (*labeled-source* (*Prule* α *As*) |- *p*) = *Some* (α, *length p*)
unfolding *labeled-source.simps lhs* **using** *label-term-increase* **by** (*metis add-0*)
with *assms*(*3*) *False* **have** *False* **by** *simp*
}
with *assms*(*2*) **obtain** *p1 p2 x* **where** *p*:*p* = *p1*@*p2* **and** *p1*:*p1* ∈ *poss* (*lhs* α) *lhs* α |- *p1* = *Var x* **and** *p2*:*p2* ∈ *poss* (($\langle$*map source As*$\rangle_\alpha$) *x*)
unfolding *source.simps* **using** *poss-subst-apply-term*[*of p lhs* α] **by** *metis*
then have *p1* ∈ *var-poss* (*lhs* α) **using** *var-poss-iff* **by** *blast*
with *p1* **obtain** *i* **where** *i*:*i* < *length As vars-term-list* (*lhs* α) !*i* = *x var-poss-list* (*lhs* α) ! *i* = *p1*
using *l2* **by** (*metis in-set-conv-nth length-var-poss-list term.inject*(*1*) *var-poss-list-sound vars-term-list-var-poss-list*)
with *p2 l1* **have** *p2-poss*:*p2* ∈ *poss* (*source* (*As*!*i*))
by (*smt* (*verit*, *del-insts*) α *case-prodD left-lin left-linear-trs-def length-map lhs-subst-var-i linear-term-var-vars-term-list nth-map*)
from *p1* **have** *labeled-source* (*Prule* α *As*) |- *p* = (($\langle$*map labeled-source As*$\rangle_\alpha$) *x*)|-*p2*
unfolding *labeled-source.simps p* **by** (*smt* (*verit*) *assms*(*2*) *eval-term.simps*(*1*) *label-term-to-term labeled-source.simps*(*3*) *labeled-source-to-term p poss-term-lab-to-term subt-at-subst subterm-poss-conv var-label-term*)
moreover from *var-rule* **have** *map* ($\langle$*map labeled-source As*$\rangle_\alpha$) (*vars-term-list* (*lhs* α)) = *map labeled-source As*
by (*metis apply-lhs-subst-var-rule l1 length-map*)
ultimately have *labeled-source* (*Prule* α *As*) |- *p* = (*labeled-source* (*As*!*i*))|-*p2*
using *i* **by** (*metis map-nth-conv*)
with *assms*(*3*) **have** *get-label* (*labeled-source* (*As* ! *i*) |- *p2*) = *Some* (β, *0*) **by** *force*
with *p2-poss i p* **show** *?thesis* **by** *blast*
qed
qed
end

If the labeled source of a proof term A has the shape $t \cdot \sigma$ where all function symbols in t are unlabeled, then A matches t with some substitution τ.

context *no-var-lhs*
begin
lemma *pterm-source-substitution*:

assumes $A \in$ *wf-pterm R*
  and *source* $A = t \cdot \sigma$ **and** *linear-term t*
  and $\forall\, p \in$ *fun-poss t*. $p \notin$ *possL A*
shows $A =$ (*to-pterm t*) $\cdot$ (*mk-subst Var* (*match-substs* (*to-pterm t*) *A*))
  using *assms* **proof**(*induct A arbitrary*:*t* σ)
  case (*1 x*)
  from *1*(*1*) **obtain** *y* **where** *y*:$t =$ *Var y*
    using *subst-apply-eq-Var* **by** (*metis source.simps*(*1*))
  have *match*:*match-substs* (*Var y*) (*Var x*) = [(*y*, *Var x*)]
    unfolding *match-substs-def vars-term-list.simps poss-list.simps* **by** *simp*
  show *?case* **unfolding** *y to-pterm.simps match*
    by *simp*
next
  case (*2 As f*)
  show *?case* **proof**(*cases t*)
    case (*Var x*)
    have *match*:*match-substs* (*Var x*) (*Pfun f As*) = [(*x*, *Pfun f As*)]
      unfolding *match-substs-def vars-term-list.simps poss-list.simps* **by** *simp*
    then show *?thesis* **unfolding** *Var to-pterm.simps match* **by** *simp*
  next
    case (*Fun g ts*)
    from *2*(*2*) **have** *f*:*f* = *g*
      unfolding *Fun* **by** *simp*
    from *2*(*2*) **have** *l*:*length ts* = *length As*
      unfolding *Fun eval-term.simps* **using** *map-eq-imp-length-eq* **by** *fastforce*
    {**fix** *i* **assume** *i*:$i <$ *length As*
      from *2*(*2*) *i* **have** *source* (*As*!*i*) = (*ts*!*i*) $\cdot\ \sigma$
        unfolding *Fun f* **by** (*smt* (*verit*, *best*) *eval-term.simps*(*2*) *l nth-map source.simps*(*2*) *term.sel*(*4*))
      moreover from *2*(*3*) *i l* **have** *lin-tsi*:*linear-term* (*ts*!*i*)
        unfolding *Fun* **by** *simp*
      moreover have ($\forall\, p\in$*fun-poss* (*ts*!*i*). $p \notin$ *possL* (*As*!*i*)) **proof**
        fix *p* **assume** $p \in$ *fun-poss* (*ts*!*i*)
        then have *i#p* $\in$ *fun-poss* (*Fun f ts*)
          using *i l* **by** *simp*
        with *2*(*4*) **have** *i#p* $\notin$ *possL* (*Pfun f As*)
          unfolding *Fun f* **by** *fastforce*
        then show $p \notin$ *possL* (*As*!*i*)
          using *i* **unfolding** *labeled-source.simps labelposs.simps* **by** *simp*
      qed
      ultimately have *IH*:*As*!*i* = *to-pterm* (*ts*!*i*) $\cdot$ *mk-subst Var* (*match-substs* (*to-pterm* (*ts*!*i*)) (*As*!*i*))
        using *2*(*1*) *i nth-mem* **by** *blast*
      have *As*!*i* = *to-pterm* (*ts*!*i*) $\cdot$ *mk-subst Var* (*match-substs* (*to-pterm t*) (*Pfun g As*)) **proof**−
        {**fix** *x* **assume** $x \in$ *vars-term* (*to-pterm* (*ts*!*i*))
        then obtain *j* **where** *j*:*vars-term-list* (*ts*!*i*) !*j* = *x j* < *length* (*vars-term-list* (*ts*!*i*))
          by (*metis in-set-conv-nth set-vars-term-list vars-to-pterm*)

then have *j'*:*j* < *length* (*map* ((|-) (*As* ! *i*)) (*var-poss-list* (*to-pterm* (*ts* ! *i*))))
by (*metis length-map length-var-poss-list var-poss-list-to-pterm*)
let *?qj*=*var-poss-list* (*to-pterm* (*ts*!*i*)) !*j*
have *map-j*:(*map* ((|-) (*As* ! *i*)) (*var-poss-list* (*to-pterm* (*ts* ! *i*))))!*j* = (*As*!*i*)|-*?qj*
using *j'* **by** *simp*
have *distinct* (*vars-term-list* (*ts*!*i*))
using *lin-tsi* **by** (*metis distinct-remdups distinct-rev linear-term-var-vars-term-list o-apply*)
then have *dist1*:*distinct* (*map fst* (*match-substs* (*to-pterm* (*ts*!*i*)) (*As*!*i*)))
unfolding *match-substs-def* **by** (*metis length-map length-var-poss-list map-fst-zip vars-to-pterm*)
have *distinct* (*vars-term-list t*)
by (*metis 2.prems*(*2*) *distinct-remdups distinct-rev linear-term-var-vars-term-list o-apply*)
then have *dist2*:*distinct* (*map fst* (*match-substs* (*to-pterm t*) (*Pfun g As*)))
unfolding *match-substs-def* **by** (*metis length-map length-var-poss-list map-fst-zip vars-to-pterm*)
have (*x*, (*As*!*i*)|-*?qj*) ∈ *set* (*match-substs* (*to-pterm* (*ts*!*i*)) (*As*!*i*))
unfolding *match-substs-def* **using** *map-j j j'* **by** (*metis* (*no-types*, *lifting*) *in-set-conv-nth length-zip min-less-iff-conj nth-zip vars-to-pterm*)
then have *sub1*:*mk-subst Var* (*match-substs* (*to-pterm* (*ts*!*i*)) (*As*!*i*)) *x* = *As*!*i* |- *?qj*
using *dist1 map-of-eq-Some-iff* **unfolding** *mk-subst-def* **by** *simp*
let *?j'*=(*sum-list* (*map* (*length* ∘ *vars-term-list*) (*take i ts*)) + *j*)
have *x2*:*vars-term-list t* ! *?j'* = *x*
unfolding *Fun vars-term-list.simps* **using** *j*(*1*) **by** (*smt* (*verit*, *best*) *concat-nth i j*(*2*) *l length-map map-map nth-map take-map*)
have *lj'*:*?j'* < *length* (*vars-term-list* (*to-pterm t*)) **unfolding** *vars-to-pterm* **unfolding** *Fun to-pterm.simps vars-term-list.simps*
using *i j*(*2*) *l concat-nth-length* **by** (*metis List.map.compositionality length-map nth-map take-map*)
then have *j'-var-poss*:*?j'* < *length* (*var-poss-list* (*to-pterm t*))
by (*metis length-var-poss-list*)
then have *lj''*:*?j'* < *length* (*map* ((|-) (*Pfun g As*)) (*var-poss-list* (*to-pterm t*)))
by (*metis length-map length-var-poss-list*)
have *var-poss-list* (*to-pterm t*) ! *?j'* = *i*#*?qj*
proof−
have *l-zip*:*i* < *length* (*zip* [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*)))
by (*simp add*: *i l*)
have *zip*:*zip* [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*)) ! *i* = (*i*, *var-poss-list* (*to-pterm* (*ts* ! *i*)))
using *nth-zip* **by** (*simp add*: *i l*)
have *map2*:*map2* (*λi*. *map* ((#) *i*)) [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*)) ! *i* ! *j* = *i*#*?qj*
unfolding *nth-map*[*OF l-zip*] *zip* **using** *j'* **by** *auto*

from *l-zip* **have** *i′′*:*i* < *length* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*)))
by *simp*
have *j′′*:*j* < *length* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*)) ! *i*)
unfolding *nth-map*[*OF l-zip*] *zip* **using** *j*(*2*) **by** (*metis case-prod-conv length-map length-var-poss-list vars-to-pterm*)
{**fix** *k* **assume** *k*:*k* < *length ts*
then have *zip*:*zip* [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*)) ! *k* = (*k*, *var-poss-list* (*to-pterm* (*ts* ! *k*)))
using *nth-zip* **by** *simp*
then have *map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*)) ! *k* =
map ((#) *k*) (*var-poss-list* (*to-pterm* (*ts* ! *k*)))
using *k* **by** *simp*
then have *length* ((*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*)))!*k*) =
length (*vars-term-list* (*ts*!*k*))
using *length-var-poss-list vars-to-pterm* **by** (*metis length-map*)
with *k* **have** (*map length* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*))))!*k* =
(*map* (*length* ∘ *vars-term-list*) *ts*) ! *k* **by** *simp*
}
moreover have *length* (*map length* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*)))) = *length ts*
by *simp*
ultimately have (*map length* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length* (*map to-pterm ts*)] (*map var-poss-list* (*map to-pterm ts*)))) =
(*map* (*length* ∘ *vars-term-list*) *ts*) **by** (*simp add*: *map-nth-eq-conv*)
then show *?thesis*
unfolding *Fun to-pterm.simps var-poss-list.simps* **using** *concat-nth*[*OF i′′ j′′*] **unfolding** *map2 take-map*[*symmetric*] **by** *simp*
qed
with *lj′′* **have** (*map* ((|-) (*Pfun g As*)) (*var-poss-list* (*to-pterm t*)))!*?j′* = *Pfun g As* |- (*i*#*?qj*)
by *force*
with *x2* **have** (*x*, *Pfun g As* |-(*i*#*?qj*)) ∈ *set* (*match-substs* (*to-pterm t*) (*Pfun g As*))
unfolding *match-substs-def* **using** *lj′ lj′′* **by** (*metis* (*no-types*, *lifting*) *in-set-conv-nth length-zip min-less-iff-conj nth-zip vars-to-pterm*)
then have *sub2*:*mk-subst Var* (*match-substs* (*to-pterm t*) (*Pfun g As*)) *x* = *Pfun g As* |- (*i*#*?qj*)
using *dist2 map-of-eq-Some-iff* **unfolding** *mk-subst-def* **by** *simp*
from *sub1 sub2* **have** *mk-subst Var* (*match-substs* (*to-pterm* (*ts*!*i*)) (*As*!*i*)) *x* = *mk-subst Var* (*match-substs* (*to-pterm t*) (*Pfun g As*)) *x*
by *simp*
}
then show *?thesis* **using** *IH*
using *term-subst-eq* **by** *force*

```
      qed
    }
    then show ?thesis
     unfolding Fun f eval-term.simps to-pterm.simps using l by (metis (mono-tags,
lifting) length-map map-nth-eq-conv)
  qed
next
  case (3 α As)
  show ?case proof(cases t)
    case (Var x)
    have match:match-substs (Var x) (Prule α As) = [(x, Prule α As)]
      unfolding match-substs-def vars-term-list.simps poss-list.simps by simp
    then show ?thesis unfolding Var to-pterm.simps match by simp
  next
    case (Fun g ts)
    from 3(1) no-var-lhs obtain f ss where lhsa:lhs α = Fun f ss
      by blast
    have [] ∈ possL (Prule α As)
     unfolding labeled-source.simps lhsa label-term.simps labelposs.simps eval-term.simps
by simp
    with 3(6) have False unfolding Fun by simp
    then show ?thesis by simp
  qed
qed

lemma unlabeled-source-to-pterm:
assumes labeled-source A = s · τ
    and linear-term s and A ∈ wf-pterm R
    and labelposs s = {}
 shows ∃ As. A = to-pterm (term-lab-to-term s) · (mk-subst Var (zip (vars-term-list
s) As)) ∧ length (vars-term-list s) = length As
  using assms proof(induct s arbitrary:A)
  case (Var x)
  let ?As =[A]
 have A = to-pterm (term-lab-to-term (Var x)) · mk-subst Var (zip (vars-term-list
(Var x)) ?As)
   unfolding term-lab-to-term.simps to-pterm.simps vars-term-list.simps zip-Cons-Cons
zip-Nil mk-subst-def by simp
  then show ?case
    by (smt (verit) length-nth-simps(1) list.size(4) vars-term-list.simps(1))
next
  case (Fun fl ts)
  from Fun(5) obtain f where f:fl = (f, None)
    by (metis empty-iff empty-pos-in-poss get-label.simps(2) get-label-imp-labelposs
prod.exhaust-sel subt-at.simps(1))
  with Fun(2) have ∃ As. A = Pfun f As ∧ length As = length ts proof(cases A)
    case (Pfun g As)
    from Fun(2) show ?thesis
      unfolding Pfun f labeled-source.simps using map-eq-imp-length-eq by auto
```

next
case (*Prule* α *As*)
from *Fun*(*4*) *no-var-lhs* **obtain** *g ss* **where** *lhs*:*lhs* α = *Fun g ss*
by (*metis Inl-inject Prule case-prodD is-FunE is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *term.distinct*(*1*) *term.sel*(*2*) *wf-pterm.simps*)
from *Fun*(*2*) **show** *?thesis*
unfolding *Prule f lhs labeled-source.simps* **by** *force*
qed *simp*
then obtain *As* **where** *as*:*A* = *Pfun f As* **and** *l*:*length As* = *length ts*
by *blast*
{**fix** *i* **assume** *i*:*i* < *length ts*
with *Fun*(*2*) **have** *labeled-source* (*As*!*i*) = (*ts*!*i*) · τ
unfolding *as* **by** (*smt* (*verit*, *best*) *eval-term.simps*(*2*) *l labeled-source.simps*(*2*) *nth-map term.inject*(*2*))
moreover from *i Fun*(*3*) **have** *linear-term* (*ts*!*i*)
by *simp*
moreover from *i Fun*(*4*) **have** *As*!*i* ∈ *wf-pterm R*
unfolding *as* **by** (*metis l fun-well-arg nth-mem*)
moreover from *i Fun*(*5*) **have** *labelposs* (*ts*!*i*) = {}
unfolding *f labelposs.simps* **by** *blast*
ultimately have ∃ *As*′. (*As*!*i*) = *to-pterm* (*term-lab-to-term* (*ts*!*i*)) · *mk-subst Var* (*zip* (*vars-term-list* (*ts*!*i*)) *As*′) ∧ *length* (*vars-term-list* (*ts*!*i*)) = *length As*′
using *Fun*(*1*) *i* **by** *force*
}
then obtain *As*′ **where** *l*′′:*length As*′ = *length ts*
and *IH*:(∀ *i* < *length ts*. (*As*!*i*) = *to-pterm* (*term-lab-to-term* (*ts*!*i*)) · *mk-subst Var* (*zip* (*vars-term-list* (*ts*!*i*)) (*As*′!*i*)) ∧ *length* (*vars-term-list* (*ts*!*i*)) = *length* (*As*′!*i*))
using *Ex-list-of-length-P*[**where** *P*=λ*As*′ *i*. *As* ! *i* = *to-pterm* (*term-lab-to-term* (*ts* ! *i*)) · *mk-subst Var* (*zip* (*vars-term-list* (*ts* ! *i*)) *As*′) ∧ *length* (*vars-term-list* (*ts*!*i*)) = *length As*′] *l* **by** *blast*
then have *l*′:*length As*′ = *length* (*map to-pterm* (*map term-lab-to-term ts*))
by *simp*
have *vars-list*:*map vars-term-list* (*map to-pterm* (*map term-lab-to-term ts*)) = *map vars-term-list ts*
by (*smt* (*verit*, *best*) *length-map map-nth-eq-conv vars-term-list-term-lab-to-term vars-to-pterm*)
have *map vars-term* (*map to-pterm* (*map term-lab-to-term ts*)) = *map vars-term ts*
using *vars-term-list-term-lab-to-term* **by** (*smt* (*verit*, *ccfv-threshold*) *length-map map-nth-eq-conv set-vars-term-list vars-to-pterm*)
then have *part*:*is-partition* (*map vars-term* (*map to-pterm* (*map term-lab-to-term ts*)))
using *Fun*(*3*) **by** (*metis linear-term.simps*(*2*))
have ∗:∀ *i* < *length ts*. *to-pterm* (*term-lab-to-term* (*ts*!*i*)) · *mk-subst Var* (*concat* (*map2 zip* (*map vars-term-list* (*map to-pterm* (*map term-lab-to-term ts*))) *As*′)) = *As*!*i*
using *mk-subst-partition-special*[*OF l*′ *part*] **unfolding** *length-map* **using** *nth-map IH*

by (*smt* (*verit*, *best*) *length-map vars-term-list-term-lab-to-term vars-to-pterm*)
from *IH* **have** $\forall$ *i* $<$ *length ts*. *length* (*vars-term-list* (*to-pterm* (*term-lab-to-term* (*ts*! *i*)))) = *length* (*As*′ ! *i*)
by (*metis vars-term-list-term-lab-to-term vars-to-pterm*)
then have *ls*:$\forall$ *i* $<$ *length ts*. *length* (*map vars-term-list* (*map to-pterm* (*map term-lab-to-term ts*)) ! *i*) = *length* (*As*′ ! *i*)
using *nth-map* **by** *simp*
then have *cc*:*concat* (*map2 zip* (*map vars-term-list* (*map to-pterm* (*map term-lab-to-term ts*))) *As*′) = *zip* (*concat* (*map vars-term-list ts*)) (*concat As*′)
unfolding *vars-list* **using** *concat-map2-zip* **by** (*metis l*′ *length-map*)
have *A* = *to-pterm* (*term-lab-to-term* (*Fun fl ts*)) · *mk-subst Var* (*zip* (*vars-term-list* (*Fun fl ts*)) (*concat As*′))
unfolding *f term-lab-to-term.simps to-pterm.simps fst-conv eval-term.simps as vars-term-list.simps cc*[*symmetric*] **using** ∗ **by** (*simp add*: *l list-eq-iff-nth-eq*)
moreover have *length* (*vars-term-list* (*Fun fl ts*)) = *length* (*concat As*′)
unfolding *vars-term-list.simps*
using *l*′′ *ls* **by** (*metis eq-length-concat-nth length-map vars-list*)
ultimately show *?case* **by** *auto*
qed
end

lemma *labels-intersect-label-term*:
assumes *term-lab-to-term A* = *t* · (*term-lab-to-term* ∘ σ)
and *linear-term t*
and *labelposs A* ∩ *labelposs* ((*label-term* α *n t*) · σ) = {}
shows ∃ *As*. *A* = *term-to-term-lab t* · (*mk-subst Var* (*zip* (*vars-term-list t*) *As*)) ∧ *length As* = *length* (*vars-term-list t*)
using *assms* **proof**(*induct t arbitrary*:*A n*)
case (*Var x*)
have *A* = *mk-subst Var* (*zip* [*x*] [*A*]) *x*
unfolding *mk-subst-def* **by** *simp*
then show *?case* **unfolding** *term-to-term-lab.simps eval-term.simps vars-term-list.simps* **by** *fastforce*
next
case (*Fun f ts*)
from *Fun*(*2*) **obtain** *lab ss* **where** *a*:*A* = *Fun* (*f*, *lab*) *ss*
using *term-lab-to-term.simps* **by** (*smt* (*verit*, *ccfv-threshold*) *eroot.cases fst-conv old.prod.exhaust eval-term.simps*(*2*) *term.distinct*(*1*) *term.sel*(*2*))
from *Fun*(*4*) **have** *lab*:*lab* = *None*
unfolding *a* **using** *insertCI* **by** *auto*
from *Fun*(*2*) **have** *l*:*length ts* = *length ss*
unfolding *a* **by** (*metis length-map eval-term.simps*(*2*) *term.sel*(*4*) *term-lab-to-term.simps*(*2*))
{**fix** *i* **assume** *i*:*i* $<$ *length ts*
with *Fun*(*2*) **have** *term-lab-to-term* (*ss*!*i*) = *ts*!*i* · (*term-lab-to-term* ∘ σ)
unfolding *a term-lab-to-term.simps eval-term.simps fst-conv* **by** (*metis l nth-map term.inject*(*2*))
moreover from *i Fun*(*3*) **have** *linear-term* (*ts*!*i*)
by *simp*
moreover have *labelposs* (*ss*!*i*) ∩ *labelposs* (*label-term* α (*n*+*1*) (*ts*!*i*) · σ) =

{}
proof−
{**fix** *q* **assume** *q1*:$q \in$ *labelposs* (*ss*!*i*) **and** *q2*:$q \in$ *labelposs* (*label-term* α (*n*+*1*) (*ts*!*i*) $\cdot$ σ)
from *q1 i l* **have** *i*#*q* $\in$ *labelposs A*
unfolding *a lab label-term.simps labelposs.simps* **by** *simp*
moreover from *q2 i* **have** *i*#*q* $\in$ *labelposs* ((*label-term* α *n* (*Fun f ts*)) $\cdot$ σ)
unfolding *label-term.simps eval-term.simps labelposs.simps length-map* **by** *simp*
ultimately have *False*
using *Fun*(*4*) **by** *blast*
}
then show *?thesis*
by *blast*
qed
ultimately have $\exists$ *As. ss*!*i* = *term-to-term-lab* (*ts*!*i*) $\cdot$ (*mk-subst Var* (*zip* (*vars-term-list* (*ts*!*i*)) *As*)) $\wedge$ *length As* = *length* (*vars-term-list* (*ts*!*i*))
using *Fun*(*1*) *i nth-mem* **by** *blast*
}
then obtain *Ass* **where** *l'*:*length Ass* = *length ts*
and *IH*:($\forall$ *i* < *length ts.* (*ss*!*i*) = (*term-to-term-lab* (*ts*!*i*)) $\cdot$ *mk-subst Var* (*zip* (*vars-term-list* (*ts*!*i*)) (*Ass*!*i*)) $\wedge$ *length* (*Ass*!*i*) = *length* (*vars-term-list* (*ts*!*i*)))
using *Ex-list-of-length-P*[**where** *P*=λ*Ass i. ss* ! *i* = (*term-to-term-lab* (*ts* ! *i*)) $\cdot$ *mk-subst Var* (*zip* (*vars-term-list* (*ts* ! *i*)) *Ass*) $\wedge$ *length Ass* = *length* (*vars-term-list* (*ts*!*i*))] *l* **by** *blast*
let *?As*=*concat Ass*
from *l'* **have** *l''*:*length Ass* = *length* (*map term-to-term-lab ts*)
by *simp*
have *vars-list*:*map vars-term-list* (*map term-to-term-lab ts*) = *map vars-term-list ts*
using *vars-term-list-term-to-term-lab* **by** *auto*
have *map vars-term* (*map term-to-term-lab ts*) = *map vars-term ts*
using *vars-term-list-term-to-term-lab* **by** (*smt* (*verit*, *ccfv-threshold*) *length-map map-nth-eq-conv set-vars-term-list vars-to-pterm*)
then have *part*:*is-partition* (*map vars-term* (*map term-to-term-lab ts*))
using *Fun*(*3*) **by** (*metis linear-term.simps*(*2*))
have $*$:$\forall$ *i* < *length ts.* (*term-to-term-lab* (*ts*!*i*)) $\cdot$ *mk-subst Var* (*concat* (*map2 zip* (*map vars-term-list* (*map term-to-term-lab ts*)) *Ass*)) = *ss*!*i*
using *mk-subst-partition-special*[*OF l'' part*] **unfolding** *length-map* **using** *nth-map IH*
by (*smt* (*verit*, *best*) *length-map vars-term-list-term-to-term-lab vars-to-pterm*)
from *IH* **have** $\forall$ *i* < *length ts. length* (*vars-term-list* (*term-to-term-lab* (*ts*! *i*))) = *length* (*Ass* ! *i*)
by (*metis vars-term-list-term-to-term-lab*)
then have *ls*:$\forall$ *i* < *length ts. length* (*map vars-term-list* (*map term-to-term-lab ts*) ! *i*) = *length* (*Ass* ! *i*)
using *nth-map* **by** *simp*
then have *cc*:*concat* (*map2 zip* (*map vars-term-list* (*map term-to-term-lab ts*)) *Ass*) = *zip* (*concat* (*map vars-term-list ts*)) (*concat Ass*)

unfolding *vars-list* **using** *concat-map2-zip* **by** (*metis l′ length-map*)
have A = *term-to-term-lab* (*Fun f ts*) · *mk-subst Var* (*zip* (*vars-term-list* (*Fun f ts*)) *?As*)
unfolding *term-to-term-lab.simps eval-term.simps vars-term-list.simps a lab cc*[*symmetric*] **using** ∗ **by** (*simp add*: *l list-eq-iff-nth-eq*)
moreover from *IH l′* **have** *l′′*:*length ?As* = *length* (*vars-term-list* (*Fun f ts*))
unfolding *vars-term-list.simps* **by** (*simp add*: *eq-length-concat-nth*)
ultimately show *?case*
by *blast*
qed

lemma *labeled-wf-pterm-rule-in-TRS*:
assumes $A \in$ *wf-pterm R* **and** $p \in$ *poss* (*labeled-source A*)
and *get-label* (*labeled-source A* |- *p*) = *Some* (α, *n*)
shows *to-rule* α ∈ *R*
using *assms* **proof**(*induct A arbitrary*: *p n*)
case (*2 ts f*)
from *2*(*2*,*3*) **obtain** *i p′* **where** *p*:*p* = *i#p′ i* < *length ts p′* ∈ *poss* (*labeled-source* (*ts*!*i*)) *get-label* (*labeled-source* (*ts*!*i*) |- *p′*) = *Some* (α, *n*)
unfolding *labeled-source.simps get-label.simps* **by** *auto*
with *2*(*1*) **show** *?case*
using *nth-mem* **by** *blast*
next
case (*3* β *As*)
from *3*(*4*) **consider** *p* ∈ *fun-poss* (*labeled-lhs* β) | (∃ *p1 p2 x*. *p* = *p1*@*p2*
∧ *p1* ∈ *poss* (*labeled-lhs* β) ∧ (*labeled-lhs* β)|-*p1* = *Var x*
∧ *p2* ∈ *poss* ((⟨*map labeled-source As*⟩$_\beta$) *x*)
∧ (*labeled-source* (*Prule* β *As*))|-*p* = ((⟨*map labeled-source As*⟩$_\beta$) *x*)|-*p2*)
unfolding *labeled-source.simps* **by** (*meson poss-is-Fun-fun-poss poss-subst-choice*)
then show *?case* **proof**(*cases*)
case *1*
then have *p* ∈ *fun-poss* (*lhs* β)
by (*simp add*: *fun-poss-label-term*)
then have *get-label* ((*labeled-source* (*Prule* β *As*))|-*p*) = *Some* (β, *length p*)
unfolding *labeled-source.simps* **by** (*simp add*: *label-term-increase*)
with *3*(*1*,*5*) **show** *?thesis* **by** *auto*
next
case *2*
then obtain *p1 p2 x* **where** *p1p2*:*p* = *p1* @ *p2* **and** *x*:*p1* ∈ *poss* (*labeled-lhs* β) ∧ *labeled-lhs* β |- *p1* = *Var x*
and *p2*:*p2* ∈ *poss* ((⟨*map labeled-source As*⟩$_\beta$) *x*)
and *lab*:*labeled-source* (*Prule* β *As*) |- *p* = (⟨*map labeled-source As*⟩$_\beta$) *x* |- *p2*
by *blast*
from *x* **have** *x* ∈ *vars-term* (*lhs* β)
by (*metis subt-at-imp-supteq subteq-Var-imp-in-vars-term vars-term-labeled-lhs*)
with *x* **obtain** *i* **where** *i*:*i* < *length* (*var-rule* β) ∧ (*var-rule* β)!*i* = *x*
by (*metis in-set-conv-nth set-vars-term-list vars-term-list-vars-distinct*)

with *3*(*2*) **have** $*$:($\langle$*map labeled-source As*$\rangle_\beta$) *x* = *labeled-source* (*As*!*i*)
  by (*metis* (*no-types*, *lifting*) *length-map lhs-subst-var-i nth-map*)
with *3*(*5*) *lab* **have** *get-label* ((*labeled-source* (*As*!*i*))|-*p2*) = *Some* (α, *n*)
  by *simp*
with *3*(*3*) *p2 i 3*(*2*) $*$ **show** *?thesis* **by** *force*
qed
qed *simp*

context *no-var-lhs*
begin
lemma *unlabeled-above-p*:
  assumes $A \in$ *wf-pterm R*
    and $p \in$ *poss* (*source A*)
    and $\forall\ r.\ r <_p p \longrightarrow r \notin$ *possL A*
  shows $p \in$ *poss A* $\wedge$ *labeled-source A*|-*p* = *labeled-source* (*A*|-*p*)
  using *assms* **proof**(*induct p arbitrary*: *A*)
  case (*Cons i p*)
  from *Cons*(*3*) **obtain** *f ts* **where** *f*:*source A* = *Fun f ts* **and** *i*:*i* < *length ts* **and** *p*:$p \in$ *poss* (*ts*!*i*)
    using *args-poss* **by** *blast*
  from *Cons*(*4*) **have** [] $\notin$ *possL A*
    by (*simp add*: *order-pos.less-le*)
  then have *no-lab*:*get-label* (*labeled-source A*) = *None*
    by (*metis empty-pos-in-poss get-label-imp-labelposs subt-at.simps*(*1*))
  from *Cons*(*3*) **obtain** *f′ As* **where** *a*:*A* = *Fun f′ As*
    by (*metis Cons-poss-Var eroot.cases source.simps*(*1*))
  then have *f′*:*f′* = *Inr f* **proof**(*cases A*)
    case (*Pfun g Bs*)
    then show *?thesis* **using** *f Pfun a* **by** *simp*
  next
    case (*Prule* α *Bs*)
    with *Cons*(*2*) **obtain** *g ss* **where** *g*:*lhs* α = *Fun g ss*
     using *no-var-lhs* **by** (*metis Inl-inject case-prodD is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *term.collapse*(*2*) *term.distinct*(*1*) *term.sel*(*2*) *wf-pterm.simps*)
    then show *?thesis* **using** *no-lab* **unfolding** *Prule* **by** *simp*
  qed *simp*
  from *i a* **have** *i′*:*i* < *length As*
    using *f f′* **by** *force*
  from *Cons*(*3*) **have** $p \in$ *poss* (*source* (*As*!*i*))
    unfolding *a f′* **by** *auto*
  moreover
  {**fix** *r* **assume** $r \in$ *poss* (*source* (*As*!*i*)) **and** *le*:$r <_p p$
    then have *i*#*r* $\in$ *poss* (*labeled-source A*)
      unfolding *a f′* **using** *i′* **by** *simp*
    moreover from *le* **have** *i*#*r* $<_p$ *i*#*p*
      by *simp*
    ultimately have *i*#*r* $\notin$ *possL A*
      using *Cons*(*4*) **by** *blast*
    then have $r \notin$ *possL* (*As*!*i*)

unfolding *a f′ labeled-source.simps* **using** *i′* **by** *force*
}
ultimately have $p \in$ *poss* (*As*!*i*) $\wedge$ *labeled-source* (*As*!*i*) |- *p* = *labeled-source* ((*As*!*i*) |- *p*)
using *Cons*(*1*,*2*) *i′* **unfolding** *a f′* **by** (*meson fun-well-arg nth-mem possL-subset-poss-source subsetD*)
with *i′ a f′* **show** *?case*
by *simp*
qed *simp*
end

lemma (**in** *single-redex*) *labeled-source-at-pq*:*labeled-source* (*A*|-*q*) = (*labeled-source A*)|-*p*
using *a pq q p a-well* **proof**(*induct q arbitrary*:*p A*)
case *Nil*
then have *p* = []
by (*simp add*: *subt-at-ctxt-of-pos-term subt-at-id-imp-eps*)
then show *?case*
by *simp*
next
case (*Cons i q*)
from *Cons*(*4*) **obtain** *fs Bs* **where** *a*:*A* = *Fun fs Bs* **and** *i*:*i* < *length Bs* **and** *q*:*q* $\in$ *poss* (*Bs*!*i*)
using *args-poss* **by** *blast*
let *?As* = *map* (λ*j*. (*Bs*!*i*) |- (*q* @ [*j*])) [*0*..<*length* (*var-rule* α)]
have (*map* (λ*ia*. *A* |- ((*i* # *q*) @ [*ia*])) [*0*..<*length* (*var-rule* α)]) = *?As*
unfolding *a* **by** *simp*
with *a i q Cons*(*2*,*4*) **have** *bsi*:*Bs*!*i* = (*ctxt-of-pos-term q* (*Bs*!*i*))⟨*Prule* α *?As*⟩
by (*metis ctxt-supt-id subt-at.simps*(*2*) *subt-at-ctxt-of-pos-term*)
from *Cons*(*6*) **have** *bi-well*:*Bs* ! *i* $\in$ *wf-pterm R*
unfolding *a* **by** (*meson fun-well-arg i nth-mem*)
show *?case* **proof**(*cases fs*)
case (*Inl* β)
from *Cons*(*6*) **have** *lin*:*linear-term* (*lhs* β)
unfolding *a Inl* **using** *left-lin left-linear-trs-def term.inject*(*2*) *wf-pterm.cases* **by** *fastforce*
from *Cons*(*6*) **have** *is-Fun*:*is-Fun* (*lhs* β)
unfolding *a Inl* **using** *no-var-lhs* **using** *wf-pterm.cases* **by** *auto*
from *Cons*(*6*) **have** *l-bs*:*length Bs* = *length* (*var-rule* β)
unfolding *a Inl* **using** *wf-pterm.cases* **by** *auto*
obtain *p1 p2* **where** *p*:*p* = *p1*@*p2* **and** *p1*:*p1* = *var-poss-list* (*lhs* β) ! *i* **and** *p2*:*p2* $\in$ *poss* (*source* (*Bs*!*i*))
using *ctxt-rule-obtain-pos Cons*(*4*,*5*,*3*) *lin l-bs* **unfolding** *a Inl* **by** *metis*
have *ctxt*:*ctxt-of-pos-term p2* (*source* (*Bs* ! *i*)) = *source-ctxt* (*ctxt-of-pos-term q* (*Bs* ! *i*))
proof−
from *p1* **have** *p1-pos*:*p1* $\in$ *poss* (*lhs* β)
using *i l-bs lin* **by** (*metis length-var-poss-list linear-term-var-vars-term-list nth-mem var-poss-imp-poss var-poss-list-sound*)

from *p1-pos* **have** *p1'*:*p1* $\in$ *poss* (*lhs* β $\cdot$ $\langle$*map source Bs*$\rangle_\beta$)
by *simp*
from *p1* **have** *p1''*:*var-poss-list* (*lhs* β) ! *length* (*take i Bs*) = *p1*
using *i* **by** *force*
have $*$:*lhs* β $\cdot$ $\langle$*map source Bs*$\rangle_\beta$ |- *p1* = *source* (*Bs*!*i*)
unfolding *p1* **using** *l-bs i*
by (*smt* (*verit*) *length-map lhs-subst-var-i lin linear-term-var-vars-term-list nth-map p1 p1-pos eval-term.simps*(*1*) *subt-at-subst vars-term-list-var-poss-list*)
from *Cons*(*3*) **show** *?thesis*
unfolding *a Inl p source.simps ctxt-of-pos-term.simps source-ctxt.simps Let-def ctxt-of-pos-term-append*[*OF p1'*] $*$ *p1''*
using *ctxt-comp-equals*[*OF p1'*] *p1-pos* **using** *poss-imp-subst-poss* **by** *blast*
qed
from *Cons*(*1*)[*OF bsi ctxt q p2 bi-well*] **have** *IH*: *labeled-source* (*Bs* ! *i* |- *q*) = *labeled-source* (*Bs* ! *i*) |- *p2*
by *presburger*
from *p1* **have** *p1* = *var-poss-list* (*labeled-lhs* β) ! *i*
by (*simp add*: *var-poss-list-labeled-lhs*)
moreover then have (*labeled-lhs* β)|-*p1* = *Var* (*vars-term-list* (*lhs* β)!*i*)
by (*metis i l-bs lin linear-term-var-vars-term-list vars-term-list-labeled-lhs vars-term-list-var-poss-list*)
ultimately show *?thesis*
unfolding *a Inl* **using** *i IH* **unfolding** *subt-at.simps p labeled-source.simps*
by (*smt* (*verit*, *ccfv-threshold*) *apply-lhs-subst-var-rule filter-cong l-bs length-map length-var-poss-list lin linear-term-var-vars-term-list map-nth-conv nth-mem poss-imp-subst-poss eval-term.simps*(*1*) *subt-at-append subt-at-subst var-poss-imp-poss var-poss-list-sound vars-term-list-labeled-lhs*)
next
case (*Inr f*)
from *Cons*(*3*,*5*) **obtain** *p'* **where** *p*:*p* = *i*#*p'* **and** *p'*:*p'* $\in$ *poss* (*source* (*Bs*!*i*))
by (*metis Cons.prems*(*3*) *Inr a source-poss*)
from *Cons*(*3*) **have** *ctxt*:*ctxt-of-pos-term p'* (*source* (*Bs* ! *i*)) = *source-ctxt* (*ctxt-of-pos-term q* (*Bs* ! *i*))
unfolding *a Inr p* **by** (*simp add*: *i*)
from *Cons*(*1*)[*OF bsi ctxt q p' bi-well*] **have** *IH*:*labeled-source* (*Bs* ! *i* |- *q*) = *labeled-source* (*Bs* ! *i*) |- *p'*
by *presburger*
then show *?thesis*
unfolding *a Inr p* **using** *i* **by** *simp*
qed
qed

context *left-lin*
begin
lemma *single-redex-label*:
assumes Δ = *ll-single-redex s p* α *p* $\in$ *poss s q* $\in$ *poss* (*source* Δ) *to-rule* α $\in$ *R*
and *get-label* (*labeled-source* Δ |-*q*) = *Some* (β, *n*)
shows $\alpha = \beta \wedge (\exists\, q'.\ q = p@q' \wedge$ *length* $q' = n \wedge q' \in$ *fun-poss* (*lhs* α))
proof –

from *assms* **have** *wf*:$\Delta \in$ *wf-pterm R*
  using *single-redex-wf-pterm left-lin left-linear-trs-def* **by** *fastforce*
from *assms* **have** $q \in$ *possL* Δ
  using *get-label-imp-labelposs* **by** *force*
then obtain q' **where** *q*:$q = p@q'$ **and** q':$q' \in$ *fun-poss* (*lhs* α)
  unfolding *assms*(*1*) **using** *single-redex-possL*[*OF assms*(*4*,*2*)] **by** *auto*
from *assms* **have** *labeled-source* Δ = (*ctxt-of-pos-term p* (*labeled-source* (*to-pterm s*)))⟨*labeled-source* (*Prule* α (*map* (*to-pterm* $\circ$ ($\lambda pi.$ *s*|-(*p*@*pi*))) (*var-poss-list* (*lhs* α)))))⟩
  using *label-source-ctxt* **by** (*simp add*: *ll-single-redex-def p-in-poss-to-pterm source-ctxt-to-pterm*)
then have *labeled-source* Δ |-*q* = *labeled-source* (*Prule* α (*map* (*to-pterm* $\circ$ ($\lambda pi.$ *s*|-(*p*@*pi*))) (*var-poss-list* (*lhs* α)))) |-q'
  unfolding *q* **using** *assms*(*2*) **by** (*metis hole-pos-ctxt-of-pos-term hole-pos-poss labeled-source-to-term poss-term-lab-to-term replace-at-subt-at source-to-pterm subt-at-append*)

then have *get-label* (*labeled-source* Δ |-*q*) = *get-label* (*labeled-lhs* α|-q')
  using *get-label-at-fun-poss-subst* q' **by** *force*
also have ... = *Some* (α, *size* q')
  using *get-label-label-term* q' **by** *fastforce*
finally show *?thesis* **using** *assms q* q' **by** *force*
qed
end

5.2 Measuring Overlap

abbreviation *measure-ov* :: ('*f*, '*v*) *pterm* $\Rightarrow$ ('*f*, '*v*) *pterm* $\Rightarrow$ *nat*
  where *measure-ov A B* $\equiv$ *card* ((*possL A*) $\cap$ (*possL B*))

lemma *finite-labelposs*: *finite* (*labelposs A*)
  by (*meson finite-fun-poss labelposs-subs-fun-poss rev-finite-subset*)

lemma *finite-possL*: *finite* (*possL A*)
  by (*simp add*: *finite-labelposs*)

lemma *measure-ov-symm*: *measure-ov A B* = *measure-ov B A*
  by (*simp add*: *Int-commute*)

lemma *measure-lhs-subst*:
  assumes *l*:*length As* = *length Bs*
  shows *card* ((*labelposs* ((*label-term* α *j t*) $\cdot$ ⟨*map labeled-source As*⟩$_\alpha$)) $\cap$
      (*labelposs* (*labeled-source* (*to-pterm t*) $\cdot$ ⟨*map labeled-source Bs*⟩$_\alpha$)))
    = ($\sum$ *x*←*vars-term-list t*. *measure-ov* ((⟨*As*⟩$_\alpha$) *x*) ((⟨*Bs*⟩$_\alpha$) *x*))
  using *assms* **proof**(*induct t arbitrary*:*j*)
case (*Var x*)
  show *?case* **proof**(*cases* $\exists$ *i* < *length As*. *i* < *length* (*var-rule* α) $\wedge$ *x* = (*var-rule* α)!*i*)
  case *True*
  then obtain *i* **where** *i*:*x* = (*var-rule* α)!*i* **and** *il*:*i* < *length As* **and** *il2*:*i* <

length (*var-rule* α) **by** *auto*
then have *a*:($\langle$*map labeled-source As*$\rangle_\alpha$) *x* = *labeled-source* (*As*!*i*)
using *lhs-subst-var-i* **by** (*metis* (*no-types*, *lifting*) *length-map nth-map*)
from *i il il2 l* **have** *b*:($\langle$*map labeled-source Bs*$\rangle_\alpha$) *x* = *labeled-source* (*Bs*!*i*)
using *lhs-subst-var-i* **by** (*metis* (*no-types*, *lifting*) *length-map nth-map*)
from *i* **show** *?thesis* **unfolding** *vars-term-list.simps sum-list-elem*
unfolding *to-pterm.simps label-term.simps labeled-source.simps eval-term.simps*
unfolding *a b* **using** *lhs-subst-var-i l il il2* **by** *metis*
next
case *False*
then have *a*:($\langle$*map labeled-source As*$\rangle_\alpha$) *x* = *Var x*
using *lhs-subst-not-var-i* **by** (*metis length-map*)
from *False l* **have** *b*:($\langle$*map labeled-source Bs*$\rangle_\alpha$) *x* = *Var x*
using *lhs-subst-not-var-i* **by** (*metis length-map*)
from *False l* **have** *possL* (($\langle$*As*$\rangle_\alpha$) *x*) $\cap$ *possL* (($\langle$*Bs*$\rangle_\alpha$) *x*) = {}
unfolding *term.set*(*3*) **using** *lhs-subst-not-var-i*
by (*metis inf.idem labeled-source.simps*(*1*) *labelposs.simps*(*1*))
then show *?thesis* **unfolding** *label-term.simps to-pterm.simps labeled-source.simps eval-term.simps a b*
by *auto*
qed
next
case (*Fun f ts*)
let *?as*=(*map* (λ*t*. *t* $\cdot$ $\langle$*map labeled-source As*$\rangle_\alpha$) (*map* (*label-term* α (*j* + *1*)) *ts*))
let *?bs*=(*map* (λ*t*. *t* $\cdot$ $\langle$*map labeled-source Bs*$\rangle_\alpha$) (*map labeled-source* (*map to-pterm ts*)))
let *?f*=(λ*i*. ({*i* # *p* |*p*. *p* $\in$ *labelposs* (*?as* ! *i*)} $\cap$ {*i* # *p* |*p*. *p* $\in$ *labelposs* (*?bs* ! *i*)}))
have {[]} $\cap$ ($\bigcup$*i*<*length ts*. {*i* # *p* |*p*. *p* $\in$ *labelposs* (*map* (λ*t*. *t* $\cdot$ $\langle$*map labeled-source Bs*$\rangle_\alpha$) (*map labeled-source* (*map to-pterm ts*)) ! *i*)}) = {}
by *blast*
then have $*$:*labelposs* (*label-term* α *j* (*Fun f ts*) $\cdot$ $\langle$*map labeled-source As*$\rangle_\alpha$) $\cap$ *labelposs* (*labeled-source* (*to-pterm* (*Fun f ts*)) $\cdot$ $\langle$*map labeled-source Bs*$\rangle_\alpha$)
= ($\bigcup$*i*<*length ts*. (*?f i*))
unfolding *label-term.simps to-pterm.simps labeled-source.simps eval-term.simps labelposs.simps* **by** *auto*
have *is-partition* (*map ?f* [*0*..<*length ts*]) **proof**−
{**fix** *i j* **assume** *j*:*j* < *length ts* **and** *i*:*i* < *j*
have *?f i* $\cap$ *?f j* = {} **unfolding** *Int-def* **using** *i*
by *fastforce*
}
then show *?thesis* **unfolding** *is-partition-def* **by** *auto*
qed
moreover have $\forall$ *i*<*length ts*. *finite* (*?f i*) **by** (*simp add*: *finite-labelposs*)
ultimately have $**$:*card* ($\bigcup$*i*<*length ts*. (*?f i*)) = ($\sum$*i*<*length ts*. *card* (*?f i*))
unfolding $*$ **using** *card-Union-Sum* **by** *blast*
{**fix** *i* **assume** *i*:*i* < *length ts*
have *?f i* = {*i* # *p* |*p*. *p* $\in$ *labelposs* (*?as* ! *i*) $\cap$ *labelposs* (*?bs* ! *i*)}
unfolding *Int-def* **by** *blast*

then have *card* (*?f i*) = *card* (*labelposs* (*?as* ! *i*) $\cap$ *labelposs* (*?bs* ! *i*))
unfolding *Setcompr-eq-image* **using** *card-image* **by** (*metis* (*no-types*, *lifting*) *inj-on-Cons1*)
with *Fun i* **have** *card* (*?f i*) = ($\sum$ *x*$\leftarrow$*vars-term-list* (*ts*!*i*). *measure-ov* (($\langle As\rangle_\alpha$) *x*) (($\langle Bs\rangle_\alpha$) *x*))
by *simp*
}
then show *?case* **unfolding** *vars-term-list.simps* $*$ $**$
by (*simp add*: *sum-sum-concat*)
qed

lemma *measure-lhs-args-zero*:
assumes *l*:*length As* = *length Bs*
and *empty*:$\forall$ *i* < *length As*. *measure-ov* (*As*!*i*) (*Bs*!*i*) = *0*
shows *measure-ov* (*Prule* α *As*) ((*to-pterm* (*lhs* α)) $\cdot$ $\langle Bs\rangle_\alpha$) = *0*
proof−
let *?xs*=*vars-term-list* (*lhs* α)
have *sum*:*measure-ov* (*Prule* α *As*) ((*to-pterm* (*lhs* α)) $\cdot$ $\langle Bs\rangle_\alpha$)
= ($\sum$ *x*$\leftarrow$*vars-term-list* (*lhs* α). *measure-ov* (($\langle As\rangle_\alpha$) *x*) (($\langle Bs\rangle_\alpha$) *x*))
using *labeled-source-apply-subst measure-lhs-subst*[*OF l*]
by (*metis* (*mono-tags*, *lifting*) *fun-mk-subst labeled-source.simps*(*1*) *labeled-source.simps*(*3*) *to-pterm-wf-pterm*)
{fix *i* **assume** *i*:*i* < *length ?xs*
have *measure-ov* (($\langle As\rangle_\alpha$) (*?xs* ! *i*)) (($\langle Bs\rangle_\alpha$) (*?xs* ! *i*)) = *0*
proof(*cases* ($\exists$*j*<*length As*. *j* < *length* (*var-rule* α) $\wedge$ (*?xs*!*i*) = *var-rule* α ! *j*))
case *True*
then obtain *j* **where** *j*:*j* < *length As j* < *length* (*var-rule* α) **and** *ij*:*?xs*!*i* = (*var-rule* α)!*j*
by *blast*
then show *?thesis*
unfolding *ij* **using** *empty* **by** (*metis j l lhs-subst-var-i*)
next
case *False*
then have ($\langle As\rangle_\alpha$) (*?xs*!*i*) = *Var* (*?xs*!*i*)
using *lhs-subst-not-var-i* **by** *metis*
moreover have ($\langle Bs\rangle_\alpha$) (*?xs*!*i*) = *Var* (*?xs*!*i*)
using *l False lhs-subst-not-var-i* **by** *metis*
ultimately show *?thesis* **by** *simp*
qed}
then show *?thesis*
using *sum* **by** (*simp add*: *sum-list-zero*)
qed

lemma *measure-zero-subt-at*:
assumes *term-lab-to-term A* = *term-lab-to-term B*
and *labelposs A* $\cap$ *labelposs B* = {}
and *p* $\in$ *poss A*
shows *labelposs* (*A*|-*p*) $\cap$ *labelposs* (*B*|-*p*) = {}
using *assms* **proof**(*induct p arbitrary*: *A B*)

case (*Cons i p*)
from *Cons*(*4*) **obtain** *f a ts* **where** *a*:$A = Fun\ (f,\ a)\ ts$ **and** *i*:$i < length\ ts$ **and** *p*:$p \in poss\ (ts!i)$
using *args-poss* **by** (*metis old.prod.exhaust*)
with *Cons*(*2*) **obtain** *b ss* **where** *b*:$B = Fun\ (f,\ b)\ ss$
by (*metis* (*no-types*, *opaque-lifting*) *Cons.prems*(*3*) *Term.term.simps*(*2*) *args-poss old.prod.exhaust poss-term-lab-to-term prod.sel*(*1*) *term-lab-to-term.simps*(*2*))
have *ts*:$(\bigcup i{<}length\ ts.\ \{i\ \#\ p\ |\ p.\ p \in labelposs\ (ts\ !\ i)\}) \subseteq labelposs\ A$ **unfolding** *a* **by**(*cases a*) *auto*
have *ss*:$(\bigcup i{<}length\ ss.\ \{i\ \#\ p\ |\ p.\ p \in labelposs\ (ss\ !\ i)\}) \subseteq labelposs\ B$ **unfolding** *b* **by**(*cases b*) *auto*
from *ss ts b i Cons*(*2*,*3*,*4*) **have** $labelposs\ (ts!i) \cap labelposs\ (ss!i) = \{\}$ **by** *auto*
with *Cons*(*1*,*2*) *p i* **show** *?case*
unfolding *a b* **by** (*simp add*: *map-eq-conv'*)
qed *simp*

lemma *empty-step-imp-measure-zero*:
assumes *is-empty-step A*
shows $measure\text{-}ov\ A\ B = 0$
by (*metis assms card-eq-0-iff inf-bot-left labeled-source-simple-pterm source-empty-step*)

lemma *measure-ov-to-pterm*:
shows $measure\text{-}ov\ A\ (to\text{-}pterm\ t) = 0$
by (*simp add*: *labeled-source-simple-pterm*)

lemma *measure-zero-imp-orthogonal*:
assumes *R*:*left-lin-no-var-lhs R* **and** *S*:*left-lin-no-var-lhs S*
and *co-initial A B* $A \in wf\text{-}pterm\ R$ $B \in wf\text{-}pterm\ S$
and $measure\text{-}ov\ A\ B = 0$
shows $A \perp_p B$
using *assms*(*3*−) **proof**(*induct A arbitrary*:*B rule*:*subterm-induct*)
case (*subterm A*)
then show *?case* **proof**(*cases A*)
case (*Var x*)
with *subterm* **show** *?thesis* **proof**(*cases B*)
case (*Prule* α *Bs*)
from *subterm*(*2*) *Var* **obtain** *y* **where** *y*:$lhs\ \alpha = Var\ y$
unfolding *Prule* **by** (*metis source.simps*(*1*) *source.simps*(*3*) *subst-apply-eq-Var*)
from *subterm*(*4*) *Prule S* **have** *is-Fun* (*lhs* α)
unfolding *left-lin-no-var-lhs-def no-var-lhs-def*
by (*metis Inl-inject case-prodD is-FunI is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *is-VarI term.inject*(*2*) *wf-pterm.simps*)
with *y* **show** *?thesis* **by** *simp*
qed (*simp-all add*: *orthogonal.intros*(*1*))
next
case (*Pfun f As*)
note *A*=*this*
with *subterm* **show** *?thesis* **proof**(*cases B*)

case (*Pfun g Bs*)
from *subterm*(*2*) **have** *f*:*f* = *g*
unfolding *Pfun A* **by** *simp*
from *subterm*(*2*) **have** *l*:*length As* = *length Bs*
unfolding *A Pfun* **using** *map-eq-imp-length-eq* **by** *auto*
{**fix** *i* **assume** *i*:*i* < *length As*
then have *As*!*i* ◁ *A*
unfolding *A* **by** *simp*
moreover from *i subterm*(*2*) *l* **have** *co-initial* (*As*!*i*) (*Bs*!*i*)
by (*metis* (*mono-tags*, *lifting*) *A Pfun nth-map source.simps*(*2*) *term.inject*(*2*))
moreover from *i subterm*(*3*) **have** *As*!*i* ∈ *wf-pterm R*
using *A* **by** *auto*
moreover from *i subterm*(*4*) *l* **have** *Bs*!*i* ∈ *wf-pterm S*
using *Pfun* **by** *auto*
moreover have *measure-ov* (*As*!*i*) (*Bs*!*i*) = *0* **proof**−
{**fix** *p* **assume** *a*:*p* ∈ *possL* (*As*!*i*) **and** *b*:*p* ∈ *possL* (*Bs*!*i*)
with *i* **have** *i*#*p* ∈ *possL A*
unfolding *A labeled-source.simps labelposs.simps* **by** *simp*
moreover from *b i l* **have** *i*#*p* ∈ *possL B*
unfolding *Pfun labeled-source.simps labelposs.simps* **by** *simp*
ultimately have *False* **using** *subterm*(*4*)
by (*metis card-gt-0-iff disjoint-iff finite-Int finite-possL less-numeral-extra*(*3*) *subterm.prems*(*4*))
}
then show *?thesis*
by (*metis card.empty disjoint-iff*)
qed
ultimately have *As*!*i* $\perp_p$ *Bs*!*i*
using *subterm*(*1*) **by** *blast*
}
then show *?thesis*
unfolding *A Pfun f* **using** *l* **by** *auto*
next
case (*Prule β Bs*)
from *subterm*(*4*) *S* **have** *lin*:*linear-term* (*lhs β*)
unfolding *Prule left-lin-no-var-lhs-def left-lin-def left-linear-trs-def* **using** *wf-pterm.cases* **by** *fastforce*
have *isfun*:*is-Fun* (*lhs β*)
using *subterm*(*4*) *S no-var-lhs.lhs-is-Fun* **unfolding** *Prule left-lin-no-var-lhs-def* **by** *blast*
have (*lhs β*) · (*term-lab-to-term* ∘ (⟨*map labeled-source Bs*⟩$_β$)) = *lhs β* · ⟨*map source Bs*⟩$_β$
by (*metis label-term-to-term labeled-source.simps*(*3*) *labeled-source-to-term source.simps*(*3*) *term-lab-to-term-subst*)
with *subterm*(*2*) **have** *co-init*:*term-lab-to-term* (*labeled-source A*) = *lhs β* · (*term-lab-to-term* ∘ ⟨*map labeled-source Bs*⟩$_β$)
unfolding *Prule* **by** *simp*
from *subterm*(*5*) **have** *possL A* ∩ *possL B* = {}
by (*simp add*: *finite-possL*)

then obtain τ **where** *labeled-source* A = *term-to-term-lab* (*lhs* β) $\cdot$ τ
unfolding *labeled-source.simps*(*3*) *Prule* **using** *labels-intersect-label-term*[*OF co-init lin*] **by** *blast*
moreover have *labelposs* (*term-to-term-lab* (*lhs* β)) = {}
using *labelposs-term-to-term-lab* **by** *blast*
moreover from *lin* **have** *linear-term* (*term-to-term-lab* (*lhs* β))
using *linear-term-to-term-lab* **by** *blast*
moreover have *term-lab-to-term* (*term-to-term-lab* (*lhs* β)) = *lhs* β
by *simp*
ultimately obtain σ **where** *sigma*:A = *to-pterm* (*lhs* β) $\cdot$ σ
using *no-var-lhs.unlabeled-source-to-pterm subterm*(*3*) *R* **unfolding** *left-lin-no-var-lhs-def* **by** *metis*
let *?As*=*map* σ (*var-rule* β)
from *sigma* **have** *a*:A = (*to-pterm* (*lhs* β)) $\cdot$ $\langle$*?As*$\rangle_\beta$
by (*smt* (*verit*, *best*) *apply-lhs-subst-var-rule comp-apply length-map map-eq-conv set-remdups set-rev set-vars-term-list term-subst-eq vars-to-pterm*)
{**fix** *i* **assume** *i*:*i* < *length* (*var-rule* β)
let *?xi*=*var-rule* β!*i*
from *i* **obtain** *i*$'$ **where** *i*$'$:*i*$'$ < *length* (*vars-term-list* (*lhs* β)) *?xi* = *vars-term-list* (*lhs* β)!*i*$'$
by (*metis comp-apply in-set-conv-nth set-remdups set-rev*)
have *l*:*length Bs* = *length* (*var-rule* β)
using *subterm*(*4*) **unfolding** *Prule* **using** *wf-pterm.cases* **by** *force*
from *i* **have** *asi*:*?As*!*i* = σ *?xi*
by *simp*
then have *?As*!*i* $\lhd$ *A*
using *a sigma subst-image-subterm i*$'$ **by** (*metis is-FunE isfun nth-mem set-vars-term-list to-pterm.simps*(*2*) *vars-to-pterm*)
moreover from *i subterm*(*2*) **have** *co-initial* (*?As*!*i*) (*Bs*!*i*)
unfolding *a Prule source.simps source-apply-subst*[*OF to-pterm-wf-pterm*[*of lhs* β]] *source-to-pterm* **using** *l*
by (*smt* (*verit*, *best*) *apply-lhs-subst-var-rule comp-def i*$'$(*1*) *i*$'$(*2*) *length-map nth-map nth-mem set-vars-term-list term-subst-eq-conv*)
moreover have *measure-ov* (*?As*!*i*) (*Bs*!*i*) = *0* **proof**−
{**fix** *p* **assume** *p*:*p* $\in$ *possL* (*?As*!*i*)
let *?pi*=*var-poss-list* (*labeled-source* (*to-pterm* (*lhs* β)))!*i*$'$
have *pi*:*?pi*=*var-poss-list* (*labeled-lhs* β) !*i*$'$
by (*simp add*: *var-poss-list-term-lab-to-term*)
have *xi*:*?xi*=*vars-term-list* (*labeled-lhs* β) !*i*$'$
by (*metis i*$'$(*2*) *vars-term-list-labeled-lhs*)
have *xi*$'$:*?xi*=*vars-term-list* (*labeled-source* (*to-pterm* (*lhs* β))) ! *i*$'$
using *vars-term-list-term-lab-to-term i*$'$(*2*) **by** (*metis labeled-source-to-term source-to-pterm*)
have *i*$'$*l*:*i*$'$ < *length* (*vars-term-list* (*labeled-lhs* β))
by (*simp add*: *i*$'$(*1*) *vars-term-list-labeled-lhs*)
have *i*$'$*l*$'$:*i*$'$ < *length* (*vars-term-list* (*labeled-source* (*to-pterm* (*lhs* β))))
by (*simp add*: *i*$'$(*1*) *vars-term-list-term-lab-to-term*)
have (*labeled-source* (*to-pterm* (*lhs* β))) |-*?pi* = *Var ?xi*
using *i*$'$ **using** *i*$'$*l*$'$ *vars-term-list-var-poss-list xi*$'$ **by** *auto*

moreover have *possL A* = *labelposs* ((*labeled-source* (*to-pterm* (*lhs* β))) $\cdot$ (*labeled-source* $\circ$ σ))
using *labeled-source-apply-subst to-pterm-wf-pterm* **unfolding** *sigma* **by** *metis*
with *p* **have** *?pi@p* $\in$ *possL A*
unfolding *set-labelposs-subst asi xi'* **using** *i'l'* **by** *fastforce*
with *subterm*(*5*) **have** *?pi@p* $\notin$ *possL B*
by (*meson card-eq-0-iff disjoint-iff finite-Int finite-labelposs*)
moreover {**assume** *p* $\in$ *possL* (*Bs*!*i*)
then have *?pi@p* $\in$ {*?pi* @ *q* |*q*. *q* $\in$ *labelposs* (($\langle$*map labeled-source Bs*$\rangle_\beta$) *?xi*)}
by (*smt* (*verit*) *Inl-inject Inr-Inl-False Prule apply-lhs-subst-var-rule i length-map map-nth-conv mem-Collect-eq subterm.prems*(*3*) *term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.cases*)
then have *?pi@p* $\in$ *possL B*
unfolding *Prule labeled-source.simps set-labelposs-subst xi pi* **using** *i'l* **by** *blast*
}
ultimately have *p* $\notin$ *possL* (*Bs*!*i*)
by *blast*
}
then have *possL* (*?As*!*i*) $\cap$ *possL* (*Bs*!*i*) = {}
by *blast*
then show *?thesis* **by** *simp*
qed
ultimately have *?As*!*i* $\perp_p$ *Bs*!*i*
using *subterm*(*1*,*3*,*4*) *i* **unfolding** *a Prule*
by (*smt* (*verit*, *best*) *Inr-Inl-False Term.term.simps*(*4*) *length-map lhs-subst-args-wf-pterm nth-mem sum.inject*(*1*) *term.inject*(*2*) *wf-pterm.simps*)
}
then show *?thesis* **unfolding** *a Prule* **using** *orthogonal.intros*(*4*)[*of ?As Bs*]
by (*smt* (*verit*, *best*) *Prule Term.term.simps*(*4*) *in-set-zip length-map old.sum.inject*(*1*) *prod.case-eq-if subterm.prems*(*3*) *sum.distinct*(*1*) *term.inject*(*2*) *wf-pterm.cases*)
qed *simp*
next
case (*Prule* α *As*)
then have *A*:*A* = *Prule* α *As*
by *simp*
from *Prule subterm*(*3*) *R* **have** *lin*:*linear-term* (*lhs* α)
unfolding *left-lin-no-var-lhs-def left-lin-def left-linear-trs-def*
using *wf-pterm.simps* **by** *fastforce*
obtain *f ts* **where** *f*:*lhs* α = *Fun f ts*
using *subterm*(*3*) *R no-var-lhs.lhs-is-Fun* **unfolding** *left-lin-no-var-lhs-def Prule* **by** *blast*
show *?thesis* **proof**(*cases B*)
case (*Var x*)
then show *?thesis*
by (*metis source.simps*(*1*) *source-orthogonal subterm.prems*(*1*) *to-pterm.simps*(*1*))

next
case (*Pfun g Bs*)
have (*lhs* α) · (*term-lab-to-term* ∘ (⟨*map labeled-source As*⟩$_\alpha$)) = *lhs* α · ⟨*map source As*⟩$_\alpha$
by (*metis label-term-to-term labeled-source.simps*(*3*) *labeled-source-to-term source.simps*(*3*) *term-lab-to-term-subst*)
with *subterm*(*2*) **have** *co-init*:*term-lab-to-term* (*labeled-source B*) = *lhs* α · (*term-lab-to-term* ∘ ⟨*map labeled-source As*⟩$_\alpha$)
unfolding *Prule* **by** *simp*
from *subterm*(*5*) **have** *possL A* ∩ *possL B* = {}
by (*simp add*: *finite-possL*)
then obtain τ **where** *labeled-source B* = *term-to-term-lab* (*lhs* α) · τ
unfolding *labeled-source.simps*(*3*) *Prule* **using** *labels-intersect-label-term*[*OF co-init lin*] **by** *blast*
moreover have *labelposs* (*term-to-term-lab* (*lhs* α)) = {}
using *labelposs-term-to-term-lab* **by** *blast*
moreover from *lin* **have** *linear-term* (*term-to-term-lab* (*lhs* α))
using *linear-term-to-term-lab* **by** *auto*
moreover have *term-lab-to-term* (*term-to-term-lab* (*lhs* α)) = *lhs* α
by *simp*
ultimately obtain σ **where** *sigma*:*B* = *to-pterm* (*lhs* α) · σ
using *no-var-lhs.unlabeled-source-to-pterm S subterm*(*4*) **unfolding** *left-lin-no-var-lhs-def* **by** *metis*
let *?Bs*=*map* σ (*var-rule* α)
from *sigma* **have** *b*:*B* = (*to-pterm* (*lhs* α)) · ⟨*?Bs*⟩$_\alpha$
by (*smt* (*verit*, *best*) *apply-lhs-subst-var-rule comp-apply length-map map-eq-conv set-remdups set-rev set-vars-term-list term-subst-eq vars-to-pterm*)
{**fix** *i* **assume** *i*:*i* < *length* (*var-rule* α)
let *?xi*=*var-rule* α!*i*
from *i* **obtain** *i*′ **where** *i*′:*i*′ < *length* (*vars-term-list* (*lhs* α)) *?xi* = *vars-term-list* (*lhs* α)!*i*′
by (*metis comp-apply in-set-conv-nth set-remdups set-rev*)
from *i* **have** *asi*:*?Bs*!*i* = σ *?xi*
by *simp*
moreover have *l*:*length As* = *length* (*var-rule* α)
using *subterm*(*3*) **unfolding** *A* **using** *wf-pterm.cases* **by** *force*
then have *As*!*i* ⊲ *A*
using *i* **unfolding** *A* **by** *simp*
moreover from *i subterm*(*2*) **have** *co-initial* (*As*!*i*) (*?Bs*!*i*)
unfolding *b Prule source.simps source-apply-subst*[*OF to-pterm-wf-pterm*[*of lhs* α]] *source-to-pterm* **using** *l*
by (*smt* (*verit*, *best*) *apply-lhs-subst-var-rule comp-def i*′(*1*) *i*′(*2*) *length-map nth-map nth-mem set-vars-term-list term-subst-eq-conv*)
moreover have *measure-ov* (*As*!*i*) (*?Bs*!*i*) = *0* **proof**−
{**fix** *p* **assume** *p*:*p* ∈ *possL* (*?Bs*!*i*)
let *?pi*=*var-poss-list* (*labeled-source* (*to-pterm* (*lhs* α)))!*i*′
have *pi*:*?pi*=*var-poss-list* (*labeled-lhs* α) !*i*′
by (*simp add*: *var-poss-list-term-lab-to-term*)
have *xi*:*?xi*=*vars-term-list* (*labeled-lhs* α) !*i*′

by (*metis* $i'(2)$ *vars-term-list-labeled-lhs*)
have xi':*?xi=vars-term-list* (*labeled-source* (*to-pterm* (*lhs* α))) ! i'
using *vars-term-list-term-lab-to-term* $i'(2)$ **by** (*metis labeled-source-to-term source-to-pterm*)
have $i'l$:$i' <$ *length* (*vars-term-list* (*labeled-lhs* α))
by (*simp add*: $i'(1)$ *vars-term-list-labeled-lhs*)
have $i'l'$:$i' <$ *length* (*vars-term-list* (*labeled-source* (*to-pterm* (*lhs* α))))
by (*simp add*: $i'(1)$ *vars-term-list-term-lab-to-term*)
have (*labeled-source* (*to-pterm* (*lhs* α))) |-*?pi* = *Var ?xi*
using i' **using** $i'l'$ *vars-term-list-var-poss-list* xi' **by** *auto*
moreover have *possL B* = *labelposs* ((*labeled-source* (*to-pterm* (*lhs* α))) $\cdot$ (*labeled-source* $\circ$ σ))
using *labeled-source-apply-subst to-pterm-wf-pterm* **unfolding** *sigma* **by** *metis*
with *p* **have** *?pi*@*p* $\in$ *possL B*
unfolding *set-labelposs-subst asi* xi' **using** $i'l'$ **by** *fastforce*
with *subterm*(*5*) **have** *?pi*@*p* $\notin$ *possL A*
by (*meson card-eq-0-iff disjoint-iff finite-Int finite-labelposs*)
moreover **{assume** *p* $\in$ *possL* (*As*!*i*)
then have *?pi*@*p* $\in$ {*?pi* @ *q* |*q*. *q* $\in$ *labelposs* (($\langle$*map labeled-source As*$\rangle_\alpha$) *?xi*)}
by (*smt* (*verit*) *Inl-inject Inr-Inl-False Prule apply-lhs-subst-var-rule i length-map map-nth-conv mem-Collect-eq subterm.prems*(*2*) *term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.cases*)
then have *?pi*@*p* $\in$ *possL A*
unfolding *Prule labeled-source.simps set-labelposs-subst xi pi* **using** $i'l$ **by** *blast*
}
ultimately have *p* $\notin$ *possL* (*As*!*i*)
by *blast*
}
then have *possL* (*As*!*i*) $\cap$ *possL* (*?Bs*!*i*) = {}
by *blast*
then show *?thesis* **by** *simp*
qed
ultimately have *As*!*i* $\perp_p$ *?Bs*!*i*
using *subterm*(*1*,*3*,*4*) *i* **unfolding** *b Prule*
by (*smt* (*verit*, *best*) *Inr-Inl-False Term.term.simps*(*4*) *length-map lhs-subst-args-wf-pterm nth-mem sum.inject*(*1*) *term.inject*(*2*) *wf-pterm.simps*)
}
then show *?thesis* **unfolding** *b Prule* **using** *orthogonal.intros*(*3*)[*of As ?Bs*]
by (*smt* (*verit*, *best*) *Prule Term.term.simps*(*4*) *in-set-zip length-map old.sum.inject*(*1*) *prod.case-eq-if subterm.prems*(*2*) *sum.distinct*(*1*) *term.inject*(*2*) *wf-pterm.cases*)
next
case (*Prule* β *Bs*)
from *subterm*(*4*) *S* **obtain** *g ss* **where** *g*:*lhs* β = *Fun g ss*
unfolding *Prule left-lin-no-var-lhs-def* **using** *no-var-lhs.lhs-is-Fun* **by** *blast*
have [] $\in$ *possL A*

unfolding *A f labeled-source.simps label-term.simps eval-term.simps label-poss.simps* **by** *blast*
moreover have [] $\in$ *possL B*
unfolding *Prule g labeled-source.simps label-term.simps eval-term.simps labelposs.simps* **by** *blast*
ultimately show *?thesis*
using *subterm*(*5*) **by** (*simp add*: *disjoint-iff finite-labelposs*)
qed
qed
qed

5.3 Collecting Overlapping Positions

abbreviation *overlaps-pos* :: ($'f$, $'v$) *term-lab* $\Rightarrow$ ($'f$, $'v$) *term-lab* $\Rightarrow$ (*pos* $\times$ *pos*) *set*
where *overlaps-pos A B* $\equiv$ *Set.filter* (λ(*p*,*q*). *get-label* (*A*|-*p*) $\neq$ *None* $\wedge$ *get-label* (*B*|-*q*) $\neq$ *None* $\wedge$
snd (*the* (*get-label* (*A*|-*p*))) = *0* $\wedge$ *snd* (*the* (*get-label* (*B*|-*q*))) = *0* $\wedge$
(*p* $<_p$ *q* $\wedge$ *get-label* (*A*|-*q*) $\neq$ *None* $\wedge$ *fst* (*the* (*get-label* (*A*|-*q*))) = *fst* (*the* (*get-label* (*A*|-*p*))) $\wedge$ *snd* (*the* (*get-label* (*A*|-*q*))) = *length* (*the* (*remove-prefix p q*)) $\vee$
(*q* $\leq_p$ *p* $\wedge$ *get-label* (*B*|-*p*) $\neq$ *None* $\wedge$ *fst* (*the* (*get-label* (*B*|-*q*))) = *fst* (*the* (*get-label* (*B*|-*p*))) $\wedge$ *snd* (*the* (*get-label* (*B*|-*p*))) = *length* (*the* (*remove-prefix q p*)))))
(*fun-poss A* $\times$ *fun-poss B*)

lemma *overlaps-pos-symmetric*:
assumes (*p*,*q*) $\in$ *overlaps-pos A B*
shows (*q*,*p*) $\in$ *overlaps-pos B A*
using *SigmaI assms less-pos-def* **by** *auto*

lemma *overlaps-pos-intro*:
assumes *q*@*q*$'$ $\in$ *fun-poss A* **and** *q* $\in$ *fun-poss B*
and *get-label* (*A*|-(*q*@*q*$'$)) = *Some* (γ, *0*)
and *get-label* (*B*|-*q*) = *Some* (β, *0*)
and *get-label* (*B*|-(*q*@*q*$'$)) = *Some* (β, *length q*$'$)
shows (*q*@*q*$'$, *q*) $\in$ *overlaps-pos A B*
using *assms* **by** *force*

Define the partial order on overlaps

definition *less-eq-overlap* :: *pos* $\times$ *pos* $\Rightarrow$ *pos* $\times$ *pos* $\Rightarrow$ *bool* (**infix** $\leq_o$ *50*)
where *p* $\leq_o$ *q* $\longleftrightarrow$ (*fst p* $\leq_p$ *fst q*) $\wedge$ (*snd p* $\leq_p$ *snd q*)

definition *less-overlap* :: *pos* $\times$ *pos* $\Rightarrow$ *pos* $\times$ *pos* $\Rightarrow$ *bool* (**infix** $<_o$ *50*)
where *p* $<_o$ *q* $\longleftrightarrow$ *p* $\leq_o$ *q* $\wedge$ *p* $\neq$ *q*

interpretation *order-overlaps*: *order less-eq-overlap less-overlap*
proof
show $\bigwedge x.\ x \leq_o x$

```
    by (simp add: less-eq-overlap-def)
  show ⋀x y z. x ≤o y ⟹ y ≤o z ⟹ x ≤o z
   by (smt (z3) less-eq-overlap-def less-overlap-def less-pos-def less-pos-def' less-pos-simps(5)
order-pos.dual-order.trans)
  show  ⋀x y. (x <o y) = strict (≤o) x y
    using less-eq-overlap-def less-overlap-def by fastforce
  thus ⋀x y. x ≤o y ⟹ y ≤o x ⟹ x = y
    by (meson less-overlap-def)
qed

lemma overlaps-pos-finite: finite (overlaps-pos A B)
  by (meson finite-SigmaI finite-filter finite-fun-poss)

lemma labeled-sources-imp-measure-not-zero:
  assumes p ∈ poss (labeled-source A) p ∈ poss (labeled-source B)
   and get-label ((labeled-source A)|-p) ≠ None ∧ get-label ((labeled-source B)|-p)
≠ None
  shows measure-ov A B > 0
  using assms
  by (metis card-gt-0-iff disjoint-iff finite-Int finite-possL get-label-imp-labelposs)

lemma measure-zero-imp-empty-overlaps:
  assumes measure-ov A B = 0 and co-init:co-initial A B
  shows overlaps-pos (labeled-source A) (labeled-source B) = {}
using assms(1) proof(rule contrapos-pp)
  {assume overlaps-pos (labeled-source A) (labeled-source B) ≠ {}
   then obtain p q where pq:(p, q) ∈ overlaps-pos (labeled-source A) (labeled-source
B)
      by (meson equals0D pred-equals-eq2)
    then have get-label ((labeled-source A)|-p) ≠ None ∧ get-label ((labeled-source
B)|-q) ≠ None
             ∧ (get-label ((labeled-source A)|-q) ≠ None ∨ get-label ((labeled-source
B)|-p) ≠ None)
      by auto
   moreover from pq have p ∈ poss (labeled-source A) and q ∈ poss (labeled-source
B)
      by (meson fun-poss-imp-poss mem-Sigma-iff member-filter)+
    ultimately show measure-ov A B ≠ 0
      using labeled-sources-imp-measure-not-zero co-init
      by (metis labeled-source-to-term less-numeral-extra(3) poss-term-lab-to-term)
  }
qed

lemma empty-overlaps-imp-measure-zero:
  assumes A ∈ wf-pterm R and B ∈ wf-pterm S
  and overlaps-pos (labeled-source A) (labeled-source B) = {}
  shows measure-ov A B = 0
  using assms(3) proof(rule contrapos-pp)
  {assume measure-ov A B ≠ 0
```

then obtain *p* **where** *p*:*p* $\in$ *possL A* $\wedge$ *p* $\in$ *possL B*
using *Int-emptyI* **by** *force*
then obtain α *n* **where** *a*:*get-label* ((*labeled-source A*)|-*p*) = *Some*(α, *n*)
using *possL-obtain-label* **by** *blast*
let *?p1*=*take* (*length p* $-$ *n*) *p*
obtain *q1* **where** *q1*:*p* = *?p1*@*q1*
by (*metis append-take-drop-id*)
from *a p assms*(*1*) **have** *alpha*:*get-label* (*labeled-source A* |- *?p1*) = *Some* (α, *0*) **and** *?p1* $\in$ *poss* (*labeled-source A*)
using *labelposs-subs-poss obtain-label-root* **by** *blast+*
then have *p1-pos*:*?p1* $\in$ *fun-poss* (*labeled-source A*)
using *get-label-imp-labelposs labelposs-subs-fun-poss* **by** *blast*
from *p* **obtain** β *m* **where** *b*:*get-label* ((*labeled-source B*)|-*p*) = *Some*(β, *m*)
using *possL-obtain-label* **by** *blast*
let *?p2*=*take* (*length p* $-$ *m*) *p*
obtain *q2* **where** *q2*:*p* = *?p2*@*q2*
by (*metis append-take-drop-id*)
from *b p assms*(*2*) **have** *beta*:*get-label* (*labeled-source B* |- *?p2*) = *Some* (β, *0*) **and** *?p2* $\in$ *poss* (*labeled-source B*)
using *labelposs-subs-poss obtain-label-root* **by** *blast+*
then have *p2-pos*:*?p2* $\in$ *fun-poss* (*labeled-source B*)
using *get-label-imp-labelposs labelposs-subs-fun-poss* **by** *blast*

then show *overlaps-pos* (*labeled-source A*) (*labeled-source B*) $\neq$ {} **proof**(*cases* *?p1* $\leq_p$ *?p2*)
case *True*
then obtain *p3* **where** *p2*:*?p2* = *?p1*@*p3*
by (*metis less-eq-pos-def*)
with *q2* **have** *p* = *?p1* @ *p3* @ *q2*
by *simp*
with *q1* **have** *p3*:*q1* = *p3*@*q2*
by (*metis same-append-eq*)
from *a alpha q1* **have** *n* = *length q1*
by (*metis* (*no-types*, *lifting*) *add-diff-cancel-left′ append-take-drop-id assms*(*1*) *label-term-max-value labelposs-subs-poss length-drop ordered-cancel-comm-monoid-diff-class.diff-add p same-append-eq subsetD*)
with *p3* **have** *n* = *length p3* + *length q2*
by *auto*
then have *get-label* ((*labeled-source A*)|-(*?p1*@*p3*)) = *Some* (α, *length p3*)
using *label-decrease*[*of ?p1*@*p3 q2 A*] *p1-pos a assms*(*1*)
by (*metis add.commute fun-poss-imp-poss fun-poss-term-lab-to-term labeled-source-to-term labelposs-subs-fun-poss-source p p2 q2*)
then have (*?p2*, *?p1*) $\in$ *overlaps-pos* (*labeled-source B*) (*labeled-source A*)
using *overlaps-pos-intro p1-pos p2-pos p2 alpha beta* **by** *simp*
then show *?thesis* **using** *overlaps-pos-symmetric* **by** *blast*
next
case *False*
with *q1 q2* **have** *?p2* $<_p$ *?p1*
by (*metis less-eq-pos-simps*(*1*) *pos-cases pos-less-eq-append-not-parallel*)

then obtain *p3* **where** *p2*:*?p1* = *?p2*@*p3*
using *less-pos-def'* **by** *blast*
with *q1* **have** *p* = *?p2* @ *p3* @ *q1*
by *simp*
with *q2* **have** *p3*:*q2* = *p3*@*q1*
by (*metis same-append-eq*)
from *b beta q2* **have** *m* = *length q2*
by (*metis* (*no-types*, *lifting*) *add-diff-cancel-left' append-take-drop-id assms*(*2*) *label-term-max-value labelposs-subs-poss length-drop ordered-cancel-comm-monoid-diff-class.diff-add p same-append-eq subsetD*)
with *p3* **have** *m* = *length p3* + *length q1*
by *auto*
then have *get-label* ((*labeled-source B*)|-(*?p2*@*p3*)) = *Some* (*β*, *length p3*)
using *label-decrease*[*of ?p2*@*p3 q1 B*] *p2-pos b assms*(*2*)
by (*metis add.commute fun-poss-imp-poss fun-poss-term-lab-to-term labeled-source-to-term labelposs-subs-fun-poss-source p p2 q1*)
then have (*?p1*, *?p2*) ∈ *overlaps-pos* (*labeled-source A*) (*labeled-source B*)
using *overlaps-pos-intro p1-pos p2-pos p2 alpha beta* **by** *simp*
then show *?thesis* **by** *blast*
qed
}
qed

lemma *obtain-overlap*:
assumes *p* ∈ *possL A p* ∈ *possL B*
and *get-label* (*labeled-source A*|-*p*) = *Some* (*γ*, *n*)
and *get-label* (*labeled-source B*|-*p*) = *Some* (*δ*, *m*)
and *n* ≤ *length p m* ≤ *length p*
and *rγ* = *take* (*length p* − *n*) *p*
and *rδ* = *take* (*length p* − *m*) *p*
and *rδ* $\leq_p$ *rγ*
and *a-well*:*A* ∈ *wf-pterm R* **and** *b-well*:*B* ∈ *wf-pterm S*
shows (*rγ*, *rδ*) ∈ *overlaps-pos* (*labeled-source A*) (*labeled-source B*)
proof−
from *assms*(*9*) **obtain** *r'* **where** *r'*:*rγ* = *rδ* @ *r'*
using *prefix-pos-diff* **by** *metis*
have *rδ* @ *r'* ∈ *fun-poss* (*labeled-source A*)
using *assms*(*1*,*7*) **unfolding** *r'* **by** (*metis append-take-drop-id fun-poss-append-poss' labelposs-subs-fun-poss subsetD*)
moreover have *rδ* ∈ *fun-poss* (*labeled-source B*)
using *assms*(*2*,*4*,*8*) **by** (*metis append-take-drop-id fun-poss-append-poss' labelposs-subs-fun-poss subsetD*)
moreover have *get-label* ((*labeled-source A*) |- (*rδ* @ *r'*)) = *Some* (*γ*, *0*)
using *assms*(*1*,*3*,*5*,*7*) *a-well* **unfolding** *r'* **using** *label-decrease*[*of take* (*length p* − *n*) *p drop* (*length p*− *n*) *p*]
by (*smt* (*verit*, *best*) *add.right-neutral add-diff-cancel-left' append-assoc append-take-drop-id labelposs-subs-poss le-add-diff-inverse2 length-drop subsetD*)
moreover have *get-label* ((*labeled-source B*) |- (*rδ*)) = *Some* (*δ*, *0*)
using *assms*(*2*,*4*,*6*,*8*) *b-well* **using** *label-decrease*[*of take* (*length p* − *m*) *p drop*

(*length p*− *m*) *p*]
  by (*smt* (*verit*, *best*) *add.right-neutral add-diff-cancel-left′ append-assoc append-take-drop-id labelposs-subs-poss le-add-diff-inverse2 length-drop subsetD*)
 moreover have *get-label* ((*labeled-source B*) |- (*rδ*@*r′*)) = *Some* (*δ*, *length r′*)
  using *assms*(*2*,*4*,*6*,*8*) *b-well* **unfolding** *r′* **using** *label-decrease*[*of take* (*length p* − *length r′*) *p drop* (*length p*− *length r′*) *p*]
  by (*smt* (*verit*, *del-insts*) *Nat.add-diff-assoc add-diff-cancel-left′ append.assoc append-take-drop-id assms*(*7*) *diff-diff-cancel diff-le-self fun-poss-imp-poss fun-poss-term-lab-to-term label-decrease labeled-source-to-term labelposs-subs-fun-poss-source le-add1 le-add-diff-inverse length-append length-take min.absorb2 r′*)
 ultimately show *?thesis* **using** *overlaps-pos-intro* **unfolding** *r′*
  by (*smt* (*verit*, *ccfv-threshold*) *append.assoc case-prodI fst-conv less-eq-pos-simps*(*1*) *mem-Sigma-iff member-filter option.distinct*(*1*) *option.sel remove-prefix-append snd-conv*)
qed

end

6 Redex Patterns

theory *Redex-Patterns*
imports
 Labels-and-Overlaps
begin

Collect all rule symbols of a proof term together with the position in its source where they appear. This is used to split a proof term into a set of single steps, whose union ($\bigsqcup$) is the whole proof term again.

The redex patterns are collected in leftmost outermost order.

fun *redex-patterns* :: (*′f*, *′v*) *pterm* ⇒ ((*′f*, *′v*) *prule* × *pos*) *list*
 where
 redex-patterns (*Var x*) = []
| *redex-patterns* (*Pfun f ss*) = *concat* (*map* (*λ* (*i*, *rps*). *map* (*λ* (*α*, *p*). (*α*, *i*#*p*)) *rps*)
  (*zip* [*0* ..< *length ss*] (*map redex-patterns ss*)))
| *redex-patterns* (*Prule α ss*) = (*α*, []) # *concat* (*map* (*λ* (*p1*, *rps*). *map* (*λ* (*α*, *p2*). (*α*, *p1*@*p2*)) *rps*)
  (*zip* (*var-poss-list* (*lhs α*)) (*map redex-patterns ss*)))

interpretation *lexord-linorder*:
 linorder ord.lexordp-eq ((<) :: *nat* ⇒ *nat* ⇒ *bool*)
  ord.lexordp ((<) :: *nat* ⇒ *nat* ⇒ *bool*)
 using *linorder.lexordp-linorder*[*OF linorder-class.linorder-axioms*] **by** *simp*

lemma *lexord-prefix-diff*:
 assumes (*ord.lexordp* ((<) :: *nat* ⇒ *nat* ⇒ *bool*)) *xs ys* **and** ¬ *prefix xs ys*
 shows (*ord.lexordp* (<)) (*xs*@*us*) (*ys*@*vs*)
using *assms* **proof**(*induct xs arbitrary*:*ys*)

case (*Cons x xs'*)
from *Cons*(*2*) obtain *y ys'* where *ys*:*ys* = *y#ys'*
by (*metis list.exhaust-sel ord.lexordp-simps*(*2*))
consider $x < y \mid x = y \wedge$ *ord.lexordp* (<) *xs' ys'*
using *Cons*(*2*) *ord.lexordp-eq.simps* unfolding *Cons ys* by *force*
then show *?case* proof(*cases*)
case *1*
then show *?thesis* unfolding *ys* by *simp*
next
case *2*
with *Cons*(*3*) have ¬ *prefix xs' ys'*
unfolding *ys* by *simp*
with *Cons*(*1*) *2* have (*ord.lexordp* (<)) (*xs'@us*) (*ys'@vs*)
by *auto*
then show *?thesis* unfolding *ys* using *2* by *simp*
qed
qed *simp*

lemma *var-poss-list-sorted*: *sorted-wrt* (*ord.lexordp* ((<) :: *nat* ⇒ *nat* ⇒ *bool*)) (*var-poss-list t*)
proof(*induct t*)
case (*Fun f ts*)
let *?poss*=(*map2* (λ*i*. *map* ((#) *i*)) [*0*..<*length ts*] (*map var-poss-list ts*))
{fix *i j* assume *i*:*i* < *length* (*var-poss-list* (*Fun f ts*)) and *j*:*j* < *i*
let *?p*=*concat ?poss* ! *i*
let *?q*=*concat ?poss* ! *j*
from *i* obtain *i' i''* where *p*:*?p* = *?poss*!*i'*!*i''* and *i'*:*i'* < *length ts* and *i''*:*i''* < *length* (*?poss*!*i'*)
and *i-sum*:*i* = *sum-list* (*map length* (*take i' ?poss*)) + *i''*
using *less-length-concat*[*OF i*[*unfolded var-poss-list.simps*]] unfolding *length-map* by *auto*
from *p* have *p2*:*?p* = *i'*# (*var-poss-list* (*ts*!*i'*) ! *i''*)
using *i' i''* by *simp*
from *j* obtain *j' j''* where *q*:*?q* = *?poss*!*j'*!*j''* and *j'*:*j'* < *length ts* and *j''*:*j''* < *length* (*?poss*!*j'*)
and *j-sum*:*j* = *sum-list* (*map length* (*take j' ?poss*)) + *j''*
using *less-length-concat i j length-map* unfolding *var-poss-list.simps*
by (*smt* (*verit*, *ccfv-threshold*) *diff-le-self length-take length-upt length-zip map-upt-len-conv order.strict-trans take-all*)
from *q* have *q2*:*?q* = *j'*# (*var-poss-list* (*ts*!*j'*) ! *j''*)
using *j' j''* by *simp*
have (*ord.lexordp* (<)) (*var-poss-list* (*Fun f ts*) ! *j*) (*var-poss-list* (*Fun f ts*) ! *i*)
proof(*cases i'* = *j'*)
case *True*
have *l*:*length* (*map2* (λ*x*. *map* ((#) *x*)) [*0*..<*length ts*] (*map var-poss-list ts*) ! *j'*) = *length* (*var-poss-list* (*ts*!*j'*))
using *j'* by *auto*
from *True j j-sum i-sum* have *j''* < *i''*
using *nat-add-left-cancel-less* by *blast*

with *Fun*(*1*)[*of ts*!*j'*] *i' i'' j''* **have** (*ord.lexordp* (<)) (*var-poss-list* (*ts*!*j'*) ! *j''*) (*var-poss-list* (*ts*!*i'*) ! *i''*)
unfolding *True l* **by** (*simp add*: *sorted-wrt-iff-nth-less*)
then have (*ord.lexordp* (<)) *?q ?p*
unfolding *p2 q2 True* **by** *simp*
then show *?thesis* **unfolding** *var-poss-list.simps* **by** *fastforce*
next
case *False*
then have *j' < i'*
using *i'' i' j' i-sum j-sum sum-list-less*[*OF j*]
by (*smt* (*verit*, *best*) *i j le-neq-implies-less length-concat linorder-le-less-linear not-add-less1 order.strict-trans take-all var-poss-list.simps*(*2*))
then have (*ord.lexordp* (<)) *?q ?p*
unfolding *p2 q2* **by** *simp*
then show *?thesis* **unfolding** *var-poss-list.simps* **by** *fastforce*
qed
}
then show *?case*
using *sorted-wrt-iff-nth-less* **by** *blast*
qed *simp*

context *left-lin-no-var-lhs*
begin

lemma *redex-patterns-sorted*:
assumes *A* ∈ *wf-pterm R*
shows *sorted-wrt* (*ord.lexordp* (<)) (*map snd* (*redex-patterns A*))
proof−
{fix *i j* **assume** *i < j j < length* (*redex-patterns A*)
with *assms* **have** (*ord.lexordp* (<)) (*snd* (*redex-patterns A* ! *i*)) (*snd* (*redex-patterns A* ! *j*))
proof(*induct A arbitrary*: *i j*)
case (*2 As f*)
let *?poss*=*map2* (λ*i*. *map* (λ(α, *p*). (α, *i* # *p*))) [*0*..<*length As*] (*map redex-patterns As*)
from *2*(*2*,*3*) **obtain** α *p1* **where** *ap*:*redex-patterns* (*Pfun f As*) ! *i* = (α, *p1*)
by (*metis surj-pair*)
from *2*(*3*) **obtain** β *p2* **where** *bp*:*redex-patterns* (*Pfun f As*) ! *j* = (β, *p2*)
by (*metis surj-pair*)
have *l*:*length* (*zip* [*0*..<*length As*] (*map redex-patterns As*)) = *length As* **by** *simp*
from *2*(*2*,*3*) **have** ∗:*i* < *length* (*concat ?poss*) **by** *simp*
from *ap* **obtain** *i' i''* **where** *ap1*:(α, *p1*) = *?poss*!*i'*!*i''* **and** *i'*:*i'* < *length As* **and** *i''*:*i''* < *length* (*?poss*!*i'*)
and *i*:*i* = *sum-list* (*map length* (*take i' ?poss*)) + *i''*
unfolding *redex-patterns.simps* **using** *less-length-concat*[*OF* ∗] **by** (*metis l length-map*)
have *poss-i'*:*?poss*!*i'* = *map* (λ(α, *p*). (α, *i'* # *p*)) (*redex-patterns* (*As*!*i'*))
using *i' nth-zip*[*of i'*] **by** *fastforce*

from *ap1 i''* **obtain** *p1'* **where** *p1'*:*p1* = *i'#p1'* (α, *p1'*) = *redex-patterns* (*As*!*i'*) ! *i''*
unfolding *poss-i'* **by** (*smt* (*z3*) *case-prod-conv length-map nth-map old.prod.inject surj-pair*)
from *bp* **obtain** *j' j''* **where** *ap2*:(β, *p2*) = *?poss*!*j'*!*j''* **and** *j'*:*j'* < *length As* **and** *j''*:*j''* < *length* (*?poss*!*j'*)
and *j*:*j* = *sum-list* (*map length* (*take j' ?poss*)) + *j''*
unfolding *redex-patterns.simps* **using** *less-length-concat*[*OF 2*(*3*)[*unfolded redex-patterns.simps*]] **by** (*metis l length-map*)
have *poss-j'*:*?poss*!*j'* = *map* (λ(α, *p*). (α, *j'* # *p*)) (*redex-patterns* (*As*!*j'*))
using *j' nth-zip*[*of j'*] **by** *fastforce*
from *ap2 j''* **obtain** *p2'* **where** *p2'*:*p2* = *j'#p2'* (β, *p2'*) = *redex-patterns* (*As*!*j'*) ! *j''*
unfolding *poss-j'* **by** (*smt* (*z3*) *case-prod-conv length-map nth-map old.prod.inject surj-pair*)
show *?case* **proof**(*cases i'* = *j'*)
case *True*
from *i j 2* **have** *i''* < *j''* **unfolding** *True* **by** *linarith*
moreover from *j''* **have** *j''* < *length* (*redex-patterns* (*As*!*j'*)) **unfolding** *poss-j'* **by** *auto*
ultimately have *ord.lexordp* (<) *p1' p2'* **using** *2*(*1*) *j' True p1'*(*2*) *p2'*(*2*) **by** (*metis nth-mem snd-eqD*)
then show *?thesis* **unfolding** *ap bp p1' p2' True* **by** *auto*
next
case *False*
with *2*(*2*) *i j* **have** *i'* < *j'* **using** *sum-list-less*[*OF 2*(*2*)] *i' j' j''*
by (*smt* (*verit*, *best*) ∗ *2.prems*(*2*) *le-neq-implies-less length-concat linorder-le-less-linear not-add-less1 redex-patterns.simps*(*2*) *take-all*)
then show *?thesis* **unfolding** *ap bp p1' p2'* **by** *fastforce*
qed
next
case (*3* γ *As*)
from *3*(*2*,*3*) **obtain** α *p1* **where** *ap*:*redex-patterns* (*Prule* γ *As*) ! *i* = (α, *p1*)
by (*metis surj-pair*)
from *3*(*3*) **obtain** β *p2* **where** *bp*:*redex-patterns* (*Prule* γ *As*) ! *j* = (β, *p2*)
by (*metis surj-pair*)
show *?case* **proof**(*cases i*)
case *0*
from *3*(*1*) *no-var-lhs* **obtain** *f ts* **where** *lhs*:*lhs* γ = *Fun f ts*
by *fastforce*
from *bp 3*(*4*) *0* **obtain** *j'* **where** *concat* (*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* γ)) (*map redex-patterns As*)) ! *j'* = (β, *p2*)
j' < *length* (*concat* (*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* γ)) (*map redex-patterns As*)))
unfolding *redex-patterns.simps* **using** *3.prems*(*2*) **by** *force*
then obtain *j1 j2* **where** *j1*:*j1* < *length* (*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* γ)) (*map redex-patterns As*))
and *j2*:*j2* < *length* (*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*)))

(*var-poss-list* (*lhs* γ)) (*map redex-patterns As*) ! *j1*)
and *j1j2*:(*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* γ)) (*map redex-patterns As*) ! *j1*) ! *j2* = (β, *p2*)
using *nth-concat-split* **by** *metis*
let *?p′*=*var-poss-list* (*lhs* γ)!*j1*
let *?rdp*=(*map redex-patterns As* ! *j1*)
from *j1* **have** *zip*:*zip* (*var-poss-list* (*lhs* γ)) (*map redex-patterns As*) ! *j1* = (*?p′*, *?rdp*)
unfolding *length-map length-zip* **using** *nth-zip* **by** *force*
with *j1j2* **have** (β, *p2*) = *map* (λ(α, *p2*). (α, *?p′* @ *p2*)) *?rdp* !*j2*
using *nth-map j1* **unfolding** *length-map* **by** *force*
moreover from *j2* **have** *j2* < *length* (*map* (λ(α, *p2*). (α, *?p′* @ *p2*)) *?rdp*)

unfolding *nth-map*[*OF j1*[*unfolded length-map*]] *zip* **by** *force*
ultimately have (β, *p2*) $\in$ *set* (*map* (λ(α, *p2*). (α, *?p′* @ *p2*)) *?rdp*)
by *simp*
moreover have *?p′* $\neq$ [] **proof**−
from *j1* **have** *?p′* $\in$ *var-poss* (*lhs* γ)
unfolding *length-map length-zip* **using** *nth-mem* **by** *fastforce*
then show *?thesis* **unfolding** *lhs var-poss.simps* **by** *force*
qed
ultimately have *p2* $\neq$ []
by *auto*
moreover from *0 ap* **have** *p1* = [] **by** *simp*
ultimately show *?thesis* **unfolding** *ap bp* **by** *simp*
next
case (*Suc n*)
let *?poss*=(*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* γ)) (*map redex-patterns As*))
from *3*(*1*,*2*) **have** *l*:*length* (*var-poss-list* (*lhs* γ)) = *length As*
using *linear-term-var-vars-term-list left-lin* **unfolding** *left-linear-trs-def* **using** *length-var-poss-list length-var-rule* **by** *auto*
from *3*(*4*,*5*) **have** $*$:*n* < *length* (*concat ?poss*) **unfolding** *Suc* **by** *simp*
from *ap* **have** *concat* (*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* γ)) (*map redex-patterns As*)) ! *n* = (α, *p1*)
unfolding *Suc* **by** *simp*
then obtain *i′ i″* **where** *ap1*:(α, *p1*) = *?poss*!*i′*!*i″* **and** *i′*:*i′* < *length ?poss* **and** *i″*:*i″* < *length* (*?poss*!*i′*)
and *n*:*n* = *sum-list* (*map length* (*take i′ ?poss*)) + *i″*
using *less-length-concat*[*OF* $*$] **by** *metis*
from *i′* **have** *i′2*:*i′* < *length* (*var-poss-list* (*lhs* γ)) **by** *simp*
obtain *p11* **where** *p11*:*?poss*!*i′* = *map* (λ(α, *p*). (α, *p11* @ *p*)) (*redex-patterns* (*As*!*i′*)) *var-poss-list* (*lhs* γ) !*i′* = *p11*
using *i′ nth-zip*[*of i′*] **by** *fastforce*
from *ap1 i″* **obtain** *p12* **where** *p12*:*p1* = *p11*@*p12* (α, *p12*) = *redex-patterns* (*As*!*i′*) ! *i″*
unfolding *p11* **by** (*smt* (*z3*) *case-prod-conv length-map nth-map old.prod.inject surj-pair*)
from *3*(*4*,*5*) *Suc* **obtain** *n′* **where** *j*:*j* = *Suc n′* **by** (*meson Suc-lessE*)

from *3*(*5*) **have** ∗:$n' <$ *length* (*concat ?poss*) **unfolding** *j* **by** *simp*
from *bp* **have** *concat* (*map2* (*λp1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* γ)) (*map redex-patterns As*)) ! $n' = (\beta,$ *p2*)
unfolding *j* **by** *simp*
then obtain j' j'' **where** *ap2*:(β, *p2*) = *?poss*!j'!j'' **and** j':$j' <$ *length ?poss* **and** j'':$j'' <$ *length* (*?poss*!j')
and n':$n' =$ *sum-list* (*map length* (*take* j' *?poss*)) + j''
using *less-length-concat*[*OF* ∗] **by** *metis*
from j' **have** j'*2*:$j' <$ *length* (*var-poss-list* (*lhs* γ)) **by** *simp*
obtain *p21* **where** *p21*:*?poss*!j' = *map* (λ(α, *p*). (α, *p21* @ *p*)) (*redex-patterns* (*As*!j')) *var-poss-list* (*lhs* γ) !j' = *p21*
using j' *nth-zip*[*of* j'] **by** *fastforce*
from *ap2* j'' **obtain** *p22* **where** *p22*:*p2* = *p21*@*p22* (β, *p22*) = *redex-patterns* (*As*!j') ! j''
unfolding *p21* **by** (*smt* (*z3*) *case-prod-conv length-map nth-map old.prod.inject surj-pair*)
show *?thesis* **proof**(*cases* $i' = j'$)
case *True*
from *n* n' *3*(*4*) **have** *ij*:$i'' < j''$ **unfolding** *True Suc j* **by** *linarith*
moreover from j'' **have** $j'' <$ *length* (*redex-patterns* (*As*!j')) **unfolding** *p21* **by** *auto*
moreover from j' *l* **have** $j' <$ *length As* **unfolding** *length-map* **by** *simp*
ultimately have *ord.lexordp* (<) *p12 p22* **using** *3*(*3*) *p22 p12* j' **unfolding** *True* **by** (*metis nth-mem snd-conv*)
with *p21*(*2*) *p11*(*2*) **show** *?thesis* **unfolding** *ap bp p22 p12 True* **by** (*simp add*: *ord.lexordp-append-leftI*)
next
case *False*
then have $i' < j'$
using *sum-list-less*[*OF 3*(*4*), **where** i'=i' **and** j'=j']
by (*smt* (*verit*) *3.prems*(*1*) *Suc Suc-less-SucD* i' *j* j' j'' *le-neq-implies-less n* n' *sum-list-less*)
then have *ord.lexordp* (<) *p11 p21*
using *p11*(*2*) *p21*(*2*) *var-poss-list-sorted*[*of lhs* γ] i'*2* j'*2* **using** *sorted-wrt-nth-less* **by** *blast*
moreover have ¬ *prefix p11 p21* **proof**−
from *False* j' i' **have** *parallel-pos* (*var-poss-list* (*lhs* γ) !i') (*var-poss-list* (*lhs* γ) !j')
unfolding *length-map length-zip* **using** *var-poss-parallel var-poss-list-sound distinct-var-poss-list*
by (*metis l min.idem nth-eq-iff-index-eq nth-mem*)
then show *?thesis* **using** *p11*(*2*) *p21*(*2*)
by (*metis less-eq-pos-simps*(*1*) *parallel-pos prefix-def*)
qed
ultimately show *?thesis* **unfolding** *ap bp p12 p22* **using** *lexord-prefix-diff* **by** *simp*
qed
qed
qed *simp*

}
then show *?thesis*
by (*metis* (*mono-tags*, *lifting*) *sorted-wrt-iff-nth-less sorted-wrt-map*)
qed

corollary *distinct-snd-rdp*:
assumes $A \in$ *wf-pterm R*
shows *distinct* (*map snd* (*redex-patterns A*))
using *redex-patterns-sorted*[*OF assms*] *lexord-linorder.strict-sorted-iff* **by** *simp*

lemma *redex-patterns-equal*:
assumes *wf*:$A \in$ *wf-pterm R*
and *sorted*:*sorted-wrt* (*ord.lexordp* (<)) (*map snd xs*) **and** *eq*:*set xs* = *set* (*redex-patterns A*)
shows *xs* = *redex-patterns A*
proof−
have *linord*:*class.linorder* (*ord.lexordp-eq* ((<) :: *nat* ⇒ *nat* ⇒ *bool*)) (*ord.lexordp* (<))
using *linorder.lexordp-linorder*[*OF linorder-class.linorder-axioms*] **by** *simp*
then have *map snd xs* = *map snd* (*redex-patterns A*)
using *linorder.strict-sorted-equal*[*OF linord redex-patterns-sorted*[*OF wf*] *sorted*] *eq* **by** *simp*
with *eq distinct-snd-rdp*[*OF wf*] **show** *?thesis*
using *distinct-map* **by** (*metis* (*mono-tags*, *opaque-lifting*) *inj-onD list.inj-map-strong*)

qed

lemma *redex-patterns-order*:
assumes $A \in$ *wf-pterm R* **and** $i < j$ **and** $j <$ *length* (*redex-patterns A*)
and *redex-patterns A* ! $i = (\alpha, p1)$ **and** *redex-patterns A* ! $j = (\beta, p2)$
shows $\neg$ $p2 \leq_p p1$
proof−
have (*ord.lexordp* (<)) *p1 p2*
using *redex-patterns-sorted*[*OF assms*(*1*)] *assms sorted-wrt-nth-less* **by** *fastforce*

then show *?thesis*
by (*metis less-eq-pos-def lexord-linorder.less-le-not-le ord.lexordp-eq-pref*)
qed

end

context *left-lin-no-var-lhs*
begin
lemma *redex-patterns-label*:
assumes $A \in$ *wf-pterm R*
shows $(\alpha, p) \in$ *set* (*redex-patterns A*) $\longleftrightarrow$ $p \in$ *poss* (*source A*) $\wedge$ *get-label* (*labeled-source A* |- *p*) = *Some* $(\alpha, 0)$
proof
{**assume** $(\alpha, p) \in$ *set* (*redex-patterns A*)

with *assms* **show** $p \in$ *poss* (*source A*) $\wedge$ *get-label* (*labeled-source A* |- *p*) = *Some* (α, *0*) **proof**(*induct arbitrary*:*p*)
case (*2 ts f*)
have *l*:*length* (*map2* (λ*i*. *map* (λ(α, *p*). (α, *i* # *p*))) [*0*..<*length ts*] (*map redex-patterns ts*)) = *length ts*
unfolding *length-map length-zip* **by** *simp*
with *2*(*2*) **obtain** *i* **where** *i*:*i* < *length ts* **and** *ap*:(α, *p*) $\in$ *set* ((*map2* (λ*i*. *map* (λ(α, *p*). (α, *i* # *p*))) [*0*..<*length ts*] (*map redex-patterns ts*))!*i*)
unfolding *redex-patterns.simps* **using** *in-set-idx* **by** (*metis nth-concat-split nth-mem*)
have (*map2* (λ*i*. *map* (λ(α, *p*). (α, *i* # *p*))) [*0*..<*length ts*] (*map redex-patterns ts*))!*i* = *map* (λ(α, *p*). (α, *i* # *p*)) (*redex-patterns* (*ts*!*i*))
using *nth-zip i* **by** *fastforce*
with *ap* **obtain** *p'* **where** *p'*:*p* = *i*#*p'* **and** (α, *p'*) $\in$ *set* (*redex-patterns* (*ts* !*i*)) **by** *auto*
with *2*(*1*) *i* **have** *p'* $\in$ *poss* (*source* (*ts*!*i*)) **and** *get-label* (*labeled-source* (*ts*!*i*)|-*p'*) = *Some* (α, *0*)
using *nth-mem* **by** *blast*+
with *i* **show** *?case* **unfolding** *p'* **by** *simp*
next
case (*3* β *As*)
from *no-var-lhs 3*(*1*) **obtain** *f ts* **where** *lhs*:*lhs* β = *Fun f ts* **by** *fastforce*
from *3*(*2*) **have** *l*:*length* (*var-poss-list* (*lhs* β)) = *length As*
using *left-lin.length-var-rule*[*OF left-lin-axioms 3*(*1*)] **by** (*simp add*: *length-var-poss-list*)

from *3*(*4*) **consider** (*root*) (α, *p*) = (β, []) | (*arg*) (α, *p*) $\in$ *set* (*concat* (*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* β)) (*map redex-patterns As*)))
unfolding *redex-patterns.simps* **by** (*meson set-ConsD*)
then show *?case* **proof**(*cases*)
case *root*
then have $\alpha = \beta$ **and** *p* = [] **by** *simp*+
then show *?thesis* **by** (*simp add*: *lhs*)
next
case *arg*
then obtain *i* **where** *i*:*i* < *length As* **and** *ap*:(α, *p*) $\in$ *set* ((*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* β)) (*map redex-patterns As*))!*i*)
using *in-set-idx l* **by** (*metis* (*no-types*, *lifting*) *length-map map-snd-zip nth-concat-split nth-mem*)
let *?p1*=(*var-poss-list* (*lhs* β))!*i*
have (*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* β)) (*map redex-patterns As*))!*i* = *map* (λ(α, *p*). (α, *?p1* @ *p*)) (*redex-patterns* (*As*!*i*))
using *nth-zip i l* **by** *fastforce*
with *ap* **obtain** *p2* **where** *p2*:*p* = *?p1*@*p2* **and** *ap2*:(α, *p2*) $\in$ *set* (*redex-patterns* (*As* !*i*)) **by** *auto*
with *3*(*3*) *i* **have** *poss*:*p2* $\in$ *poss* (*source* (*As*!*i*)) **and** *label*:*get-label* (*labeled-source* (*As*!*i*)|-*p2*) = *Some* (α, *0*)
using *nth-mem* **by** *blast*+
have *p1-poss*:*?p1* $\in$ *poss* (*lhs* β) **using** *i l*

by (*metis nth-mem var-poss-imp-poss var-poss-list-sound*)
then have *1*:$p \in$ *poss* (*source* (*Prule* β *As*))
using *poss 3*(*2*) *i l* **unfolding** *source.simps p2*
by (*smt* (*verit*, *ccfv-SIG*) *append-eq-append-conv2 comp-apply length-map length-remdups-eq length-rev length-var-poss-list nth-map poss-append-poss poss-imp-subst-poss rev-swap var-rule-pos-subst vars-term-list-var-poss-list*)
have *labeled-source* (*Prule* β *As*) |-*p* = *labeled-source* (*As*!*i*) |-*p2* **proof**−
have ($\langle$*map labeled-source As*$\rangle_\beta$) (*var-rule* β ! *i*) = *labeled-source* (*As*!*i*)
using *i 3*(*2*) **by** (*metis length-map lhs-subst-var-i nth-map*)
moreover have *labeled-lhs* β |- *?p1* = *Var* (*var-rule* β ! *i*)
using *3*(*1*) *i l* **by** (*metis case-prodD left-lin left-linear-trs-def length-var-poss-list linear-term-var-vars-term-list p1-poss var-label-term vars-term-list-var-poss-list*)
ultimately show *?thesis* **unfolding** *p2 labeled-source.simps*
by (*smt* (*verit*, *best*) *eval-term.simps*(*1*) *label-term-to-term p1-poss poss-imp-subst-poss poss-term-lab-to-term subt-at-append subt-at-subst*)
qed
with *label* **have** *2*:*get-label* (*labeled-source* (*Prule* β *As*)|-*p*) = *Some* (α, *0*)
by *presburger*
from *1 2* **show** *?thesis* **by** *simp*
qed
qed *simp*
}
{assume $p \in$ *poss* (*source A*) $\wedge$ *get-label* (*labeled-source A* |- *p*) = *Some* (α, *0*)
with *assms* **show** (α, *p*) $\in$ *set* (*redex-patterns A*) **proof**(*induct arbitrary*:*p*)
case (*2 ts f*)
from *2*(*2*) **have** $p \neq$ [] **unfolding** *labeled-source.simps* **by** *auto*
with *2*(*2*) **obtain** *i p'* **where** *p'*:*p* = *i*#*p'* $p' \in$ *poss* (*source* (*ts*!*i*)) **and** *i*:*i* < *length ts*
unfolding *source.simps* **by** *fastforce*
with *2*(*2*) **have** *get-label* (*labeled-source* (*ts*!*i*) |- *p'*) = *Some* (α, *0*)
unfolding *p' labeled-source.simps* **by** *auto*
with *2*(*1*) *i p'* **have** *IH*:(α, *p'*) $\in$ *set* (*redex-patterns* (*ts*!*i*))
using *nth-mem* **by** *blast*
from *i* **have** *i-zip*:*i* < *length* (*zip* [*0*..<*length ts*] (*map redex-patterns ts*)) **by** *simp*
from *i* **have** *zip* [*0*..<*length ts*] (*map redex-patterns ts*) ! *i* = (*i*, *redex-patterns* (*ts*!*i*))
using *nth-zip* **by** *simp*
then have (*map2* (λx. *map* (λ(α, *p*). (α, *x* # *p*))) [*0*..<*length ts*] (*map redex-patterns ts*)) ! *i* = *map* (λ(α, *p*). (α, *i* # *p*)) (*redex-patterns* (*ts*!*i*))
unfolding *nth-map*[*OF i-zip*] **by** *simp*
with *p'*(*2*) *IH* **have** (α, *p*) $\in$ *set* ((*map2* (λi. *map* (λ(α, *p*). (α, *i* # *p*))) [*0*..<*length ts*] (*map redex-patterns ts*))!*i*)
unfolding *p'* **by** *auto*
with *i-zip* **show** *?case* **using** *i* **unfolding** *redex-patterns.simps set-concat* **by** (*metis* (*no-types*, *lifting*) *UN-I length-map nth-mem*)
next
case (*3* β *As*)
with *get-label-Prule* **consider** (*1*)*p* = [] $\wedge$ $\alpha = \beta$ | ($\exists$ *p1 p2 i*. *p* = *p1* @ *p2*

∧ *i* < *length As* ∧ *var-poss-list* (*lhs* β) ! *i* = *p1*
∧ *p2* ∈ *poss* (*source* (*As* ! *i*)) ∧ *get-label* (*labeled-source* (*As* ! *i*) |- *p2*) = *Some* (α, *0*))
by (*metis wf-pterm.intros*(*3*))
then show *?case* **proof**(*cases*)
case *1*
then show *?thesis* **unfolding** *redex-patterns.simps* **by** *simp*
next
case *2*
from *3*(*1*,*2*) *left-lin* **have** *l*:*length* (*var-poss-list* (*lhs* β)) = *length As*
using *length-var-poss-list length-var-rule* **by** *auto*
from *2* **obtain** *p1 p2 i* **where** *p*:*p* = *p1* @ *p2* **and** *i*:*i* < *length As* **and** *p1*:*var-poss-list* (*lhs* β) ! *i* = *p1*
and *p2*:*p2* ∈ *poss* (*source* (*As* ! *i*)) **and** *lab*:*get-label* (*labeled-source* (*As* ! *i*) |- *p2*) = *Some* (α, *0*)
by *blast*
from *i l* **have** *i'*:*i* < *length* (*zip* (*var-poss-list* (*lhs* β)) (*map redex-patterns As*)) **by** *simp*
from *i p2 lab 3*(*3*) **have** (α, *p2*) ∈ *set* (*redex-patterns* (*As*!*i*)) **using** *nth-mem* **by** *blast*
then have (α, *p*) ∈ *set* (*map* (λ(α, *p2*). (α, *p1* @ *p2*)) (*redex-patterns* (*As*!*i*))) **using** *p* **by** *force*
then have (α, *p*) ∈ *set* ((*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* β)) (*map redex-patterns As*))!*i*)
unfolding *nth-map*[*OF i'*] *p* **using** *p1* **by** (*simp add*: *i l*)
then have (α, *p*) ∈ *set* (*concat* (*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* β)) (*map redex-patterns As*)))
unfolding *set-concat* **by** (*metis* (*no-types*, *lifting*) *UN-I i' length-map nth-mem*)
then show *?thesis* **unfolding** *redex-patterns.simps* **by** (*meson list.set-intros*(*2*))
qed
qed *simp*
}
qed

lemma *redex-pattern-rule-symbol*:
assumes *A* ∈ *wf-pterm R* (α, *p*) ∈ *set* (*redex-patterns A*)
shows *to-rule* α ∈ *R*
proof−
from *redex-patterns-label*[*OF assms*(*1*)] **have** *p* ∈ *poss* (*source A*) **and** *get-label* (*labeled-source A* |- *p*) = *Some* (α, *0*)
using *assms*(*2*) **by** *simp*+
then show *?thesis*
using *assms*(*1*) *labeled-wf-pterm-rule-in-TRS* **by** *fastforce*
qed

lemma *redex-patterns-subset-possL*:
assumes *A* ∈ *wf-pterm R*
shows *set* (*map snd* (*redex-patterns A*)) ⊆ *possL A*

using *redex-patterns-label*[*OF assms*]
by (*smt* (*verit*) *get-label-imp-labelposs imageE labeled-source-to-term list.set-map option.simps*(*3*) *poss-term-lab-to-term prod.collapse subsetI*)
end

lemma *redex-poss-empty-imp-empty-step*:
assumes *redex-patterns A* = []
shows *is-empty-step A*
using *assms* **proof**(*induct A*)
case (*Pfun f As*)
{**fix** *i* **assume** *i*:*i* < *length As*
then have *i-zip*:*i* < *length* (*zip* [*0*..<*length As*] (*map redex-patterns As*)) **by** *simp*
{**fix** *x xs* **assume** *redex-patterns* (*As*!*i*) = *x#xs*
with *i* **have** *zip* [*0*..<*length As*] (*map redex-patterns As*) ! *i* = (*i*, *x#xs*) **by** *simp*
then have (*map2* (λ*i*. *map* (λ(α, *p*). (α, *i* # *p*))) [*0*..<*length As*] (*map redex-patterns As*))!*i* ≠ []
using *nth-map i-zip* **by** *simp*
with *Pfun*(*2*) **have** *False* **unfolding** *redex-patterns.simps* **using** *i-zip concat-nth-length*
by (*metis* (*no-types*, *lifting*) *length-0-conv length-greater-0-conv length-map less-nat-zero-code*)
}
then have *redex-patterns* (*As*!*i*) = []
by (*meson list.exhaust*)
with *Pfun*(*1*) *i* **have** *is-empty-step* (*As*!*i*)
by *simp*
}
then show *?case*
by (*simp add*: *list-all-length*)
qed *simp-all*

lemma *overlap-imp-redex-poss*:
assumes *A* ∈ *wf-pterm R B* ∈ *wf-pterm R*
and *measure-ov A B* ≠ *0*
shows *redex-patterns A* ≠ []
proof
assume *redex-patterns A* = []
then have *is-empty-step A*
by (*simp add*: *redex-poss-empty-imp-empty-step*)
with *assms*(*3*) **show** *False*
by (*simp add*: *empty-step-imp-measure-zero*)
qed

lemma *redex-patterns-to-pterm*:
shows *redex-patterns* (*to-pterm s*) = []
proof(*induct s*)
case (*Fun f ts*)

```
  have l:length (map2 (λi. map (λ(α, p). (α, i # p))) [0..<length (map to-pterm
ts)] (map redex-patterns (map to-pterm ts))) = length ts
    by simp
  {fix i assume i < length ts
   with Fun have (map2 (λi. map (λ(α, p). (α, i # p))) [0..<length (map to-pterm
ts)] (map redex-patterns (map to-pterm ts)))!i = []
      by simp
  }
  with l show ?case unfolding to-pterm.simps redex-patterns.simps
   by (metis length-greater-0-conv length-nth-simps(1) less-nat-zero-code nth-concat-split)
qed simp


lemma redex-patterns-elem-fun:
  assumes (α, p) ∈ set (redex-patterns (Pfun f As))
  shows ∃ i p'. i < length As ∧ p = i#p' ∧ (α, p') ∈ set (redex-patterns (As!i))
proof−
  let ?xs=map2 (λi. map (λ(α, p). (α, i # p))) [0..<length As] (map redex-patterns
As)
   from assms obtain k where k:k < length (redex-patterns (Pfun f As)) re-
dex-patterns (Pfun f As) ! k = (α, p)
    by (metis in-set-idx)
  then obtain i j where i < length ?xs and j:j < length (?xs!i) ?xs ! i ! j = (α,
p)
    using nth-concat-split[OF k(1)[unfolded redex-patterns.simps]] by force
  then have i:i < length As by auto
   then have zip [0..<length As] (map redex-patterns As) !i = (i, redex-patterns
(As!i))
    using nth-zip by auto
  then have ?xs!i = map (λ(α, p). (α, i#p)) (redex-patterns (As!i)) using nth-map
i by auto
  with j obtain p' where p = i#p' and (α, p') ∈ set (redex-patterns (As!i))
      by (smt (verit, ccfv-threshold) case-prod-beta fst-conv imageE list.set-map
nth-mem prod.collapse snd-conv)
  with i show ?thesis by simp
qed

lemma redex-patterns-elem-rule:
  assumes (α, p) ∈ set (redex-patterns (Prule β As))
  shows p = [] ∨ (∃ i p'. i < length As ∧ i < length (var-poss-list (lhs β))
      ∧ p = (var-poss-list (lhs β)!i)@p' ∧ (α, p') ∈ set (redex-patterns (As!i)))
proof−
  let ?xs=map2 (λp1. map (λ(α, p2). (α, p1 @ p2))) (var-poss-list (lhs β)) (map
redex-patterns As)
   from assms obtain k where k:k < length (redex-patterns (Prule β As)) re-
dex-patterns (Prule β As) ! k = (α, p)
    by (metis in-set-idx)
  show ?thesis proof(cases p = [])
    case False
```

with *k* **have** $k \neq 0$
unfolding *redex-patterns.simps* **by** (*metis nth-Cons-0 prod.inject*)
with *k* **obtain** *i j* **where** $i < length\ ?xs$ **and** $j{:}j < length\ (?xs!i)\ ?xs\ !\ i\ !\ j = (\alpha,\ p)$
using *nth-concat-split less-Suc-eq-0-disj* **unfolding** *redex-patterns.simps* **by** *force*
then have $i{:}i < length\ As\ i < length\ (var\text{-}poss\text{-}list\ (lhs\ \beta))$ **by** *auto*
let $?p{=}(var\text{-}poss\text{-}list\ (lhs\ \beta))!i$
from *i* **have** $zip\ (var\text{-}poss\text{-}list\ (lhs\ \beta))\ (map\ redex\text{-}patterns\ As)\ !i = (?p,\ redex\text{-}patterns\ (As!i))$
using *nth-zip* **by** *auto*
then have $?xs!i = map\ (\lambda(\alpha,\ p).\ (\alpha,\ ?p@p))\ (redex\text{-}patterns\ (As!i))$ **using** *nth-map i* **by** *auto*
with *j* **obtain** p' **where** $p = ?p@p'$ **and** $(\alpha,\ p') \in set\ (redex\text{-}patterns\ (As!i))$
by (*smt* (*verit*, *ccfv-threshold*) *case-prod-beta fst-conv imageE list.set-map nth-mem prod.collapse snd-conv*)
with *i* **show** *?thesis* **by** *blast*
qed *simp*
qed

lemma *redex-patterns-elem-fun'*:
assumes $(\alpha,\ p) \in set\ (redex\text{-}patterns\ (As!i))$ **and** $i{:}i < length\ As$
shows $(\alpha,\ i\#p) \in set\ (redex\text{-}patterns\ (Pfun\ f\ As))$
proof−
let $?xs{=}map2\ (\lambda i.\ map\ (\lambda(\alpha,\ p).\ (\alpha,\ i\ \#\ p)))\ [0..{<}length\ As]\ (map\ redex\text{-}patterns\ As)$
from *i* **have** $zip\ [0..{<}length\ As]\ (map\ redex\text{-}patterns\ As)\ !i = (i,\ redex\text{-}patterns\ (As!i))$
using *nth-zip* **by** *auto*
then have $?xs!i = map\ (\lambda(\alpha,\ p).\ (\alpha,\ i\#p))\ (redex\text{-}patterns\ (As!i))$ **using** *nth-map i* **by** *auto*
with *assms* **have** $(\alpha,\ i\#p) \in set\ (?xs!i)$ **by** *fastforce*
moreover from *i* **have** $i < length\ ?xs$ **by** *simp*
ultimately have $*{:}(\alpha,\ i\#p) \in set\ (concat\ ?xs)$
unfolding *set-concat* **by** (*meson UN-iff nth-mem*)
then show *?thesis* **by** *simp*
qed

lemma *redex-patterns-elem-rule'*:
assumes $(\beta,\ p) \in set\ (redex\text{-}patterns\ (As!i))$ **and** $i{:}i < length\ As\ i < length\ (var\text{-}poss\text{-}list\ (lhs\ \alpha))$
shows $(\beta,\ (var\text{-}poss\text{-}list\ (lhs\ \alpha)\ !\ i)@p) \in set\ (redex\text{-}patterns\ (Prule\ \alpha\ As))$
proof−
let $?xs{=}map2\ (\lambda p1.\ map\ (\lambda(\alpha,\ p2).\ (\alpha,\ p1\ @\ p2)))\ (var\text{-}poss\text{-}list\ (lhs\ \alpha))\ (map\ redex\text{-}patterns\ As)$
let $?p{=}var\text{-}poss\text{-}list\ (lhs\ \alpha)\ !\ i$
from *i* **have** $zip\ (var\text{-}poss\text{-}list\ (lhs\ \alpha))\ (map\ redex\text{-}patterns\ As)\ !i = (?p,\ redex\text{-}patterns\ (As!i))$
using *nth-zip* **by** *auto*

then have *?xs*!*i* = *map* (λ(α, *p*). (α, *?p*@*p*)) (*redex-patterns* (*As*!*i*)) **using** *nth-map i* **by** *auto*
with *assms* **have** (β, *?p*@*p*) ∈ *set* (*?xs*!*i*) **by** *fastforce*
moreover from *i* **have** *i* < *length ?xs* **by** *simp*
ultimately have ∗:(β, *?p*@*p*) ∈ *set* (*concat ?xs*)
unfolding *set-concat* **by** (*meson UN-iff nth-mem*)
then show *?thesis* **by** *simp*
qed

lemma *redex-patterns-elem-subst*:
assumes (α, *p*) ∈ *set* (*redex-patterns* ((*to-pterm t*) · σ))
shows ∃ *p1 p2 x*. *p* = *p1*@*p2* ∧ (α, *p2*) ∈ *set* (*redex-patterns* (σ *x*)) ∧ *p1* ∈ *var-poss t* ∧ *t*|-*p1* = *Var x*
using *assms* **proof**(*induct t arbitrary*: *p*)
case (*Var x*)
then show *?case* **unfolding** *to-pterm.simps eval-term.simps* **by** *force*
next
case (*Fun f ts*)
from *Fun*(*2*) **obtain** *j* **where** *j*:*j* < *length* (*redex-patterns* (*to-pterm* (*Fun f ts*) · σ)) (*redex-patterns* (*to-pterm* (*Fun f ts*) · σ))!*j* = (α, *p*)
by (*metis in-set-idx*)
from *j* **obtain** *i k* **where** *i*:*i* < *length ts*
and *k*:*k* < *length* (*map* (λ(α, *p*). (α, *i* # *p*)) (*redex-patterns* (*to-pterm* (*ts*!*i*) · σ)))
and *rdp*:(*map* (λ(α, *p*). (α, *i* # *p*)) (*redex-patterns* (*to-pterm* (*ts*!*i*) · σ)))!*k* = (α, *p*)
using *nth-concat-split* **unfolding** *length-map to-pterm.simps eval-term.simps redex-patterns.simps* **by** *force*
from *rdp k* **obtain** *p*′ **where** *p*:*p* = *i*#*p*′
by (*smt* (*verit*, *del-insts*) *case-prod-conv list.sel*(*3*) *map-eq-imp-length-eq map-ident nth-map prod.inject surj-pair*)
from *k rdp* **have** (α, *p*′) ∈ *set* (*redex-patterns* (*to-pterm* (*ts*!*i*) · σ))
unfolding *p* **by** (*smt* (*verit*, *del-insts*) *case-prod-conv list.sel*(*3*) *map-eq-imp-length-eq map-ident nth-map nth-mem prod.inject surj-pair*)
with *Fun*(*1*) *i* **obtain** *p1 p2 x* **where** *p*′:*p*′ = *p1*@*p2* **and** *rdp2*:(α, *p2*) ∈ *set* (*redex-patterns* (σ *x*)) **and** *p1* ∈ *var-poss* (*ts*!*i*) **and** (*ts*!*i*)|-*p1* = *Var x*
by (*meson nth-mem*)
with *i* **have** *i*#*p1* ∈ *var-poss* (*Fun f ts*) *Fun f ts* |- (*i*#*p1*) = *Var x*
by *auto*
with *p*′ *rdp2* **show** *?case*
unfolding *p* **by** (*meson Cons-eq-appendI*)
qed

context *left-lin-no-var-lhs*
begin

lemma *redex-patterns-rule*′′:
assumes *rdp*:(β, *p* @ *q*) ∈ *set* (*redex-patterns* (*Prule* α *As*))
and *wf*:*Prule* α *As* ∈ *wf-pterm R*

and *p*:*p* = *var-poss-list* (*lhs* α)!*i*
and *i*:*i* < *length As*
shows (β, *q*) $\in$ *set* (*redex-patterns* (*As*!*i*))
proof−
from *wf* **obtain** *f ts* **where** *lhs*:*lhs* α = *Fun f ts*
by (*metis Inl-inject case-prodD is-FunE is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *no-var-lhs term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.simps*)
from *wf i* **have** *l*:*length As* = *length* (*var-poss-list* (*lhs* α))
by (*metis Inl-inject is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *length-var-poss-list length-var-rule term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.simps*)
with *i p* **have** *p* $\in$ *var-poss* (*Fun f ts*)
by (*metis lhs nth-mem var-poss-list-sound*)
then have *p* $\neq$ [] **by** *force*
then obtain *j p*′ **where** *j*:*j* < *length As* **and** *p*′:*p*@*q* = *var-poss-list* (*lhs* α) ! *j* @ *p*′ (β, *p*′) $\in$ *set* (*redex-patterns* (*As*!*j*))
using *redex-patterns-elem-rule*[*OF rdp*] **by** *blast*
{**assume** *j* $\neq$ *i*
then have *p* $\perp$ *var-poss-list* (*lhs* α) ! *j*
unfolding *p* **using** *i j* **by** (*metis distinct-var-poss-list l nth-eq-iff-index-eq nth-mem var-poss-list-sound var-poss-parallel*)
with *p*′(*1*) **have** *False*
by (*metis less-eq-pos-simps*(*1*) *pos-less-eq-append-not-parallel*)
}
with *p*′(*1*) *p* **have** *j* = *i* **and** *p*′ = *q* **by** *fastforce*+
with *p*′(*2*) **show** *?thesis* **by** *simp*
qed

lemma *redex-patterns-elem-subst*′:
assumes (α, *p2*) $\in$ *set* (*redex-patterns* (σ *x*)) **and** *p1*:*p1* $\in$ *poss t t*|-*p1* = *Var x*
shows (α, *p1*@*p2*) $\in$ *set* (*redex-patterns* ((*to-pterm t*) $\cdot$ σ))
using *assms* **proof**(*induct t arbitrary*: *p1*)
case (*Var x*)
then show *?case* **unfolding** *to-pterm.simps eval-term.simps* **by** *force*
next
case (*Fun f ts*)
from *Fun*(*3*,*4*) **obtain** *i p1*′ **where** *i*:*i* < *length ts* **and** *p1*:*p1* = *i*#*p1*′ **and** *p1*′:*p1*′ $\in$ *poss* (*ts*!*i*) (*ts*!*i*)|-*p1*′ = *Var x*
by *auto*
with *Fun*(*1*,*2*) **have** (α, *p1*′ @ *p2*) $\in$ *set* (*redex-patterns* (*to-pterm* (*ts*!*i*) $\cdot$ σ))
using *nth-mem* **by** *blast*
then obtain *j* **where** *j*:*j* < *length* (*redex-patterns* (*to-pterm* (*ts*!*i*) $\cdot$ σ)) *redex-patterns* (*to-pterm* (*ts*!*i*) $\cdot$ σ)!*j* = (α, *p1*′ @ *p2*)
by (*metis in-set-idx*)
let *?xs*=*map2* (λ*i*. *map* (λ(α, *p*). (α, *i* # *p*))) [*0*..<*length* (*map* (λ*s*. *s* $\cdot$ σ) (*map to-pterm ts*))] (*map redex-patterns* (*map* (λ*s*. *s* $\cdot$ σ) (*map to-pterm ts*)))
from *i j* **have** *rdp*:*?xs*!*i*!*j* = (α, *p1*@*p2*)
unfolding *p1* **by** *auto*
let *?i*=*sum-list* (*map length* (*take i ?xs*)) + *j*
from *rdp i j*(*1*) **have** (*redex-patterns* ((*to-pterm* (*Fun f ts*)) $\cdot$ σ)) ! *?i* = (α,

p1@*p2*)
  using *concat-nth*[*of i ?xs j*] **unfolding** *length-map* **by** *force*
 moreover from *i j*(*1*) **have** *?i* < *length* (*redex-patterns* (*to-pterm* (*Fun f ts*) · σ))
  using *concat-nth-length*[*of i ?xs j*] **unfolding** *length-map* **by** *force*
 ultimately show *?case*
  using *nth-mem* **by** *fastforce*
qed

lemma *redex-patterns-join*:
 assumes *A* ∈ *wf-pterm R B* ∈ *wf-pterm R A* ⊔ *B* = *Some C*
 shows *set* (*redex-patterns C*) = *set* (*redex-patterns A*) ∪ *set* (*redex-patterns B*)
 using *assms* **proof**(*induct A arbitrary*: *B C rule*:*subterm-induct*)
 case (*subterm A*)
 from *subterm*(*2*) **show** *?case* **proof**(*cases A*)
  case (*1 x*)
  from *subterm*(*2*,*3*,*4*) *var-join* **show** *?thesis*
   unfolding *1* **by** *auto*
 next
  case (*2 As f*)
  with *subterm*(*4*) **consider** (*Pfun*) ∃ *g Bs*. *B* = *Pfun g Bs* | (*Prule*) ∃ α *Bs*. *B* = *Prule* α *Bs* **by** (*meson fun-join*)
  then show *?thesis* **proof**(*cases*)
   case *Pfun*
   then obtain *g Bs* **where** *B*:*B* = *Pfun g Bs* **by** *blast*
   from *subterm*(*4*) *join-fun-fun* **obtain** *Cs* **where** *fg*:*f* = *g* **and** *l-As-Bs*:*length As* = *length Bs* **and**
     C:*C* = *Pfun f Cs* **and** *l-Cs-As*:*length Cs* = *length As* **and** *Cs*:(∀ *i*<*length As*. *As* ! *i* ⊔ *Bs* ! *i* = *Some* (*Cs* ! *i*))
    unfolding *2 B* **by** *force*
   {**fix** *i* **assume** *i*:*i* < *length As*
    with *subterm*(*3*) **have** *Bs*!*i* ∈ *wf-pterm R*
     using *B l-As-Bs* **by** *auto*
   with *subterm*(*1*) *i 2 Cs* **have** *set* (*redex-patterns* (*Cs*!*i*)) = *set* (*redex-patterns* (*As*!*i*)) ∪ *set* (*redex-patterns* (*Bs*!*i*))
     by (*meson nth-mem supt.arg*)
   }**note** *IH*=*this*
   {**fix** α *p* **assume** (α, *p*) ∈ *set* (*redex-patterns C*)
    then obtain *i p′* **where** *i*:*i* < *length Cs* **and** *p*:*p* = *i*#*p′* **and** (α, *p′*) ∈ *set* (*redex-patterns* (*Cs*!*i*))
     unfolding *C* **by** (*meson redex-patterns-elem-fun*)
     with *IH* **consider** (α, *p′*) ∈ *set* (*redex-patterns* (*As*!*i*)) | (α, *p′*) ∈ *set* (*redex-patterns* (*Bs*!*i*))
     using *l-Cs-As* **by** *fastforce*
      then have (α, *p*) ∈ *set* (*redex-patterns A*) ∪ *set* (*redex-patterns B*) **proof**(*cases*)
     case *1*
     with *i* **have** (α, *p*) ∈ *set* (*redex-patterns A*)
      unfolding *2 p l-Cs-As* **by** (*meson redex-patterns-elem-fun′*)

then show *?thesis* **by** *simp*
next
case *2*
with *i* **have** $(\alpha, p) \in$ *set* (*redex-patterns B*)
unfolding *B l-Cs-As l-As-Bs p* **by** (*meson redex-patterns-elem-fun′*)
then show *?thesis* **by** *simp*
qed
}
moreover
{fix α *p* **assume** $(\alpha, p) \in$ *set* (*redex-patterns A*) $\cup$ *set* (*redex-patterns B*)
then consider $(\alpha, p) \in$ *set* (*redex-patterns A*) | $(\alpha, p) \in$ *set* (*redex-patterns B*) **by** *force*
then have $(\alpha, p) \in$ *set* (*redex-patterns C*) **proof**(*cases*)
case *1*
then obtain *i p′* **where** *i*:*i* < *length As* **and** *p*:*p* = *i*#*p′* **and** $(\alpha, p') \in$ *set* (*redex-patterns* (*As*!*i*))
unfolding *2* **by** (*meson redex-patterns-elem-fun*)
with *IH* **have** $(\alpha, p') \in$ *set* (*redex-patterns* (*Cs*!*i*)) **by** *blast*
with *i l-Cs-As* **show** *?thesis* **unfolding** *C p*
by (*metis redex-patterns-elem-fun′*)
next
case *2*
then obtain *i p′* **where** *i*:*i* < *length Bs* **and** *p*:*p* = *i*#*p′* **and** $(\alpha, p') \in$ *set* (*redex-patterns* (*Bs*!*i*))
unfolding *B* **by** (*meson redex-patterns-elem-fun*)
with *IH l-As-Bs* **have** $(\alpha, p') \in$ *set* (*redex-patterns* (*Cs*!*i*)) **by** *simp*
with *i l-Cs-As l-As-Bs* **show** *?thesis* **unfolding** *C p*
by (*metis redex-patterns-elem-fun′*)
qed
}
ultimately show *?thesis* **by** *auto*
next
case *Prule*
then obtain α *Bs* **where** *B*:*B* = *Prule* α *Bs* **by** *blast*
from *B subterm*(*3*) **have** *alpha*:*to-rule* $\alpha \in R$
using *wf-pterm.simps* **by** *fastforce*
then obtain *f′ ts* **where** *lhs*:*lhs* α = *Fun f′ ts*
using *no-var-lhs* **by** *fastforce*
from *alpha* **have** *lin*:*linear-term* (*lhs* α)
using *left-lin left-linear-trs-def* **by** *fastforce*
from *B subterm*(*3*,*4*) **obtain** σ *Cs* **where** *sigma*:*match A* (*to-pterm* (*lhs* α)) = *Some* σ
and *C*:*C* = *Prule* α *Cs* **and** *l-Cs-Bs*:*length Cs* = *length Bs* **and** *Cs*:($\forall$ *i*<*length Bs*. σ (*var-rule* α ! *i*) $\sqcup$ (*Bs* ! *i*) = *Some* (*Cs* ! *i*))
unfolding *2* **using** *join-rule-fun join-sym* **by** (*smt* (*verit*, *best*))
from *B subterm*(*3*) **have** *l-Bs*:*length Bs* = *length* (*var-rule* α)
using *wf-pterm.simps* **by** *fastforce*
from *sigma* **have** *A*:*A* = (*to-pterm* (*lhs* α) $\cdot$ σ)
by (*simp add*: *match-matches*)

{fix *i* **assume** *i*:*i* < *length Bs*
with *sigma lhs l-Bs* **have** σ (*var-rule* α ! *i*) $\lhd$ *A*
by (*smt* (*verit*, *best*) *comp-def match-lhs-subst nth-mem set-remdups set-rev set-vars-term-list subst-image-subterm to-pterm.simps*(*2*) *vars-to-pterm*)
moreover have σ (*var-rule* α ! *i*) $\in$ *wf-pterm R*
using *subterm*(*2*) **by** (*metis i l-Bs match-well-def sigma vars-to-pterm*)
moreover from *i subterm*(*3*) **have** *Bs*!*i* $\in$ *wf-pterm R*
using *B nth-mem* **by** *blast*
ultimately have *set* (*redex-patterns* (*Cs*!*i*)) = *set* (*redex-patterns* (σ (*var-rule* α ! *i*))) $\cup$ *set* (*redex-patterns* (*Bs*!*i*))
using *subterm*(*1*) *Cs l-Cs-Bs i* **by** *presburger*
}note *IH*=*this*
{fix β *p* **assume** *rdp*:(β, *p*) $\in$ *set* (*redex-patterns C*)
then have (β, *p*) $\in$ *set* (*redex-patterns A*) $\cup$ *set* (*redex-patterns B*) **proof**(*cases p*=[])
case *True*
with *rdp* **have** $\alpha = \beta$
unfolding *C* **using** *lhs* **by** (*metis* (*no-types*, *lifting*) *C join-wf-pterm list.set-intros*(*1*) *option.sel*
prod.inject redex-patterns.simps(*3*) *redex-patterns-label subterm.prems*(*1*) *subterm.prems*(*2*) *subterm.prems*(*3*))
then show *?thesis* **unfolding** *B redex-patterns.simps True* **by** *simp*
next
case *False*
with *rdp* **obtain** *i p*′ **where** *i*:*i* < *length Cs i* < *length* (*var-poss-list* (*lhs* α))
and *p*:*p* = (*var-poss-list* (*lhs* α) ! *i*)@*p*′ **and** ∗:(β, *p*′) $\in$ *set* (*redex-patterns* (*Cs*!*i*))
unfolding *C* **by** (*meson redex-patterns-elem-rule*)
let *?p*=*var-poss-list* (*lhs* α) ! *i*
from ∗ *i IH* **consider** (β, *p*′) $\in$ *set* (*redex-patterns* (σ (*var-rule* α ! *i*))) | (β, *p*′) $\in$ *set* (*redex-patterns* (*Bs*!*i*))
using *l-Cs-Bs* **by** *fastforce*
then show *?thesis* **proof**(*cases*)
case *1*
let *?x*=*var-rule* α ! *i*
from *i*(*2*) **have** *p-pos*:*?p* $\in$ *poss* (*lhs* α)
by (*metis nth-mem var-poss-iff var-poss-list-sound*)
from *i*(*2*) **have** *p-x*:(*lhs* α)|-*?p* = *Var ?x*
by (*metis* ‹*to-rule* α $\in$ *R*› *case-prodD left-lin left-linear-trs-def length-var-poss-list linear-term-var-vars-term-list vars-term-list-var-poss-list*)
from *i*(*2*) **have** (β, *p*) $\in$ *set* (*redex-patterns A*)
unfolding *p A* **using** *redex-patterns-elem-subst*′[*of* β *p*′ σ *?x*, *OF 1 p-pos p-x*] **by** *simp*
then show *?thesis* **by** *simp*
next
case *2*
from *i* **have** (β, *p*) $\in$ *set* (*redex-patterns B*)
unfolding *B p l-Cs-Bs* **using** *redex-patterns-elem-rule*′[*OF 2*] **by**

presburger
then show *?thesis* **by** *simp*
qed
qed
}
moreover
{fix β *p* **assume** (β, *p*) $\in$ *set* (*redex-patterns A*) $\cup$ *set* (*redex-patterns B*)
then consider (β, *p*) $\in$ *set* (*redex-patterns A*) | (β, *p*) $\in$ *set* (*redex-patterns B*) **by** *force*
then have (β, *p*) $\in$ *set* (*redex-patterns C*) **proof**(*cases*)
case *1*
then obtain *p1 p2 x* **where** *p*:*p* = *p1*@*p2* **and** *rdp2*:(β, *p2*) $\in$ *set* (*redex-patterns* (σ *x*))
and *p1*:*p1* $\in$ *var-poss* (*lhs* α) *lhs* α|-*p1* = *Var x*
unfolding *A* **using** *redex-patterns-elem-subst* **by** *metis*
then obtain *i* **where** *i*:*i* < *length* (*var-rule* α) (*var-rule* α)!*i* = *x*
using *lin* **by** (*metis in-set-conv-nth length-var-poss-list linear-term-var-vars-term-list term.inject*(*1*) *var-poss-list-sound vars-term-list-var-poss-list*)
with *p1 lin* **have** *p1*:*p1* = *var-poss-list* (*lhs* α) ! *i*
by (*metis length-var-poss-list linear-term-unique-vars linear-term-var-vars-term-list nth-mem var-poss-imp-poss var-poss-list-sound vars-term-list-var-poss-list*)
from *i IH rdp2* **have** (β, *p2*) $\in$ *set* (*redex-patterns* (*Cs*!*i*))
by (*simp add*: *l-Bs*)
with *i*(*1*) **show** *?thesis* **unfolding** *C p*
using *redex-patterns-elem-rule′ p1* **by** (*metis alpha l-Bs l-Cs-Bs length-var-poss-list length-var-rule*)
next
case *2*
show *?thesis* **proof**(*cases p*=[])
case *True*
from *2* **have** $\alpha = \beta$
unfolding *B True* **using** *lhs* **by** (*metis* (*no-types*, *lifting*) *B list.set-intros*(*1*) *option.sel*
prod.inject redex-patterns.simps(*3*) *redex-patterns-label subterm.prems*(*2*))

then show *?thesis* **unfolding** *C redex-patterns.simps True* **by** *simp*
next
case *False*
with *2* **obtain** *i p′* **where** *i*:*i* < *length Bs i* < *length* (*var-poss-list* (*lhs* α))
and *p*:*p* = (*var-poss-list* (*lhs* α) ! *i*)@*p′* **and** $*$:(β, *p′*) $\in$ *set* (*redex-patterns* (*Bs*!*i*))
unfolding *B* **by** (*meson redex-patterns-elem-rule*)
with *IH l-Cs-Bs* **have** (β, *p′*) $\in$ *set* (*redex-patterns* (*Cs*!*i*)) **by** *simp*
with *i l-Cs-Bs* **show** *?thesis* **unfolding** *C p*
by (*metis redex-patterns-elem-rule′*)
qed
qed
}

ultimately show *?thesis* **by** *auto*
qed
next
case (*3* α *As*)
from *3*(*2*) **obtain** f' *ts* **where** *lhs*:*lhs* α = *Fun* f' *ts*
using *no-var-lhs* **by** *fastforce*
from *3*(*2*) **have** *lin*:*linear-term* (*lhs* α)
using *left-lin left-linear-trs-def* **by** *fastforce*
from *3 subterm*(*2*,*4*) **consider** (*Pfun*) $\exists$ *g Bs*. *B* = *Pfun g Bs* | (*Prule*) $\exists$ β *Bs*. *B* = *Prule* β *Bs* **by** (*meson rule-join*)
then show *?thesis* **proof**(*cases*)
case *Pfun*
then obtain *f Bs* **where** *B*:*B* = *Pfun f Bs* **by** *blast*
from *subterm*(*2*,*4*) **obtain** σ *Cs* **where** *sigma*:*match B* (*to-pterm* (*lhs* α)) = *Some* σ
and *C*:*C* = *Prule* α *Cs* **and** *l-Cs-As*:*length Cs* = *length As* **and** *As*:($\forall$ *i*<*length As*. (*As* ! *i*) $\sqcup$ σ (*var-rule* α ! *i*) = *Some* (*Cs* ! *i*))
unfolding *3 B* **using** *3*(*3*) *join-rule-fun* **by** *metis*
{**fix** *i* **assume** *i*:*i* < *length As*
have σ (*var-rule* α ! *i*) $\in$ *wf-pterm R*
using *subterm*(*3*) **by** (*metis 3*(*3*) *i match-well-def sigma vars-to-pterm*)
then have *set* (*redex-patterns* (*Cs*!*i*)) = *set* (*redex-patterns* (*As*!*i*)) $\cup$ *set* (*redex-patterns* (σ (*var-rule* α ! *i*)))
using *subterm*(*1*) *3* **by** (*meson As i nth-mem supt.arg*)
}**note** *IH*=*this*
{**fix** β *p* **assume** *rdp*:(β, *p*) $\in$ *set* (*redex-patterns C*)
then have (β, *p*) $\in$ *set* (*redex-patterns A*) $\cup$ *set* (*redex-patterns B*) **proof**(*cases p*=[])
case *True*
with *rdp* **have** α = β
unfolding *C* **using** *lhs* **by** (*metis* (*no-types*, *lifting*) *C join-wf-pterm list.set-intros*(*1*) *option.sel*
prod.inject redex-patterns.simps(*3*) *redex-patterns-label subterm.prems*(*1*) *subterm.prems*(*2*) *subterm.prems*(*3*))
then show *?thesis* **unfolding** *3 redex-patterns.simps True* **by** *simp*
next
case *False*
with *rdp* **obtain** *i* p' **where** *i*:*i* < *length Cs i* < *length* (*var-poss-list* (*lhs* α))
and *p*:*p* = (*var-poss-list* (*lhs* α) ! *i*)@p' **and** $*$:(β, p') $\in$ *set* (*redex-patterns* (*Cs*!*i*))
unfolding *C* **by** (*meson redex-patterns-elem-rule*)
let *?p*=*var-poss-list* (*lhs* α) ! *i*
from $*$ *i IH* **consider** (β, p') $\in$ *set* (*redex-patterns* (σ (*var-rule* α ! *i*))) | (β, p') $\in$ *set* (*redex-patterns* (*As*!*i*))
using *l-Cs-As* **by** *auto*
then show *?thesis* **proof**(*cases*)
case *1*
let *?x*=*var-rule* α ! *i*

from $i(2)$ **have** *p-pos*: $?p \in$ *poss* (*lhs* α)
by (*metis nth-mem var-poss-iff var-poss-list-sound*)
from $i(2)$ **have** *p-x*:(*lhs* α)|-$?p$ = *Var* $?x$
by (*metis* $3(2)$ *case-prodD left-lin left-linear-trs-def length-var-poss-list linear-term-var-vars-term-list vars-term-list-var-poss-list*)
from *sigma* **have** $(\beta, p) \in$ *set* (*redex-patterns B*)
unfolding *p* **using** *redex-patterns-elem-subst'*[*of* β p' σ $?x$, *OF 1 p-pos p-x*] **by** (*simp add*: *match-matches*)
then show *?thesis* **by** *simp*
next
case *2*
from *i* **have** $(\beta, p) \in$ *set* (*redex-patterns A*)
unfolding $3(1)$ *p l-Cs-As* **using** *redex-patterns-elem-rule'*[*OF 2*] **by** *presburger*
then show *?thesis* **by** *simp*
qed
qed
}
moreover
{fix β *p* **assume** $(\beta, p) \in$ *set* (*redex-patterns A*) $\cup$ *set* (*redex-patterns B*)
then consider $(\beta, p) \in$ *set* (*redex-patterns B*) | $(\beta, p) \in$ *set* (*redex-patterns A*) **by** *force*
then have $(\beta, p) \in$ *set* (*redex-patterns C*) **proof**(*cases*)
case *1*
then obtain *p1 p2 x* **where** *p*:*p* = *p1*@*p2* **and** *rdp2*:$(\beta, p2) \in$ *set* (*redex-patterns* (σ *x*))
and *p1*:*p1* $\in$ *var-poss* (*lhs* α) *lhs* α|-*p1* = *Var x*
using *sigma redex-patterns-elem-subst* **using** *match-matches* **by** *blast*
then obtain *i* **where** *i*:$i <$ *length* (*var-rule* α) (*var-rule* α)!*i* = *x*
using *lin* **by** (*metis in-set-conv-nth length-var-poss-list linear-term-var-vars-term-list term.inject*(*1*) *var-poss-list-sound vars-term-list-var-poss-list*)
with *p1 lin* **have** *p1*:*p1* = *var-poss-list* (*lhs* α) ! *i*
by (*metis length-var-poss-list linear-term-unique-vars linear-term-var-vars-term-list nth-mem var-poss-imp-poss var-poss-list-sound vars-term-list-var-poss-list*)
from *i IH rdp2* **have** $(\beta, p2) \in$ *set* (*redex-patterns* (*Cs*!*i*))
by (*simp add*: $3(3)$)
with $i(1)$ **show** *?thesis* **unfolding** *C p*
using *redex-patterns-elem-rule' p1* **by** (*metis* $3(2)$ $3(3)$ *l-Cs-As length-var-poss-list length-var-rule*)
next
case *2*
show *?thesis* **proof**(*cases p*=[])
case *True*
from *2* **have** $\alpha = \beta$
unfolding *True* **using** *lhs* **by** (*metis* $3(1)$ *list.set-intros*(*1*) *option.sel prod.sel*(*1*) *redex-patterns.simps*(*3*) *redex-patterns-label subterm.prems*(*1*))
then show *?thesis* **unfolding** *C redex-patterns.simps True* **by** *simp*
next
case *False*

with *2* **obtain** $i\ p'$ **where** i:$i <$ *length As* $i <$ *length* (*var-poss-list* (*lhs* α))
and p:$p =$ (*var-poss-list* (*lhs* α) ! i)@p' **and** $*$:(β, p') $\in$ *set* (*redex-patterns* (*As*!i))
using *3*(*1*) *redex-patterns-elem-rule* **by** *blast*
with *IH l-Cs-As* **have** (β, p') $\in$ *set* (*redex-patterns* (*Cs*!i)) **by** *simp*
with *i l-Cs-As* **show** *?thesis* **unfolding** *C p*
by (*metis redex-patterns-elem-rule'*)
qed
qed
}
ultimately show *?thesis* **by** *auto*
next
case *Prule*
then obtain β *Bs* **where** *B*:$B =$ *Prule* β *Bs* **by** *blast*
obtain *Cs* **where** *alpha-beta*:$\alpha = \beta$ **and** *l-As-Bs*:*length As* = *length Bs*
and *C*:$C =$ *Prule* α *Cs* **and** *l-Cs-As*:*length Cs* = *length As* **and** *args*:$\forall\, i <$ *length As. As* ! $i \sqcup$ *Bs* ! i = *Some* (*Cs* ! i)
using *join-rule-rule*[*OF subterm*(*4*,*2*,*3*)[*unfolded B 3*]] **using** *3*(*3*) **by** *fastforce*
{fix *i* **assume** $i <$ *length As*
from *subterm*(*1*) **have** *set* (*redex-patterns* (*Cs*!i)) = *set* (*redex-patterns* (*As*!i)) $\cup$ *set* (*redex-patterns* (*Bs*!i))
by (*metis 3*(*1*) *3*(*4*) *B* ‹$i <$ *length As*› *fun-well-arg l-As-Bs local.args nth-mem subterm.prems*(*2*) *supt.arg*)
}note *IH*=*this*
{fix γ *p* **assume** *rdp*:(γ, p) $\in$ *set* (*redex-patterns C*)
have (γ, p) $\in$ *set* (*redex-patterns A*) $\cup$ *set* (*redex-patterns B*) **proof**(*cases* $p = []$)
case *True*
from *rdp* **have** $\alpha = \gamma$ **unfolding** *C True* **using** *lhs*
by (*metis* (*no-types*, *lifting*) *C join-wf-pterm list.set-intros*(*1*) *option.sel prod.inject redex-patterns.simps*(*3*) *redex-patterns-label subterm*(*2*,*3*,*4*))
then show *?thesis* **unfolding** *3 True* **by** *simp*
next
case *False*
then obtain *p2 i* **where** i:$i <$ *length Cs* $i <$ *length* (*var-poss-list* (*lhs* α))
and p:$p =$ *var-poss-list* (*lhs* α) ! i @*p2* **and** (γ, *p2*) $\in$ *set* (*redex-patterns* (*Cs*!i))
using *C rdp redex-patterns-elem-rule* **by** *blast*
with *IH* **consider** (γ, *p2*) $\in$ *set* (*redex-patterns* (*As*!i)) | (γ, *p2*) $\in$ *set* (*redex-patterns* (*Bs*!i))
using *l-Cs-As* **by** *fastforce*
then show *?thesis* **proof**(*cases*)
case *1*
with *i* **have** (γ, p) $\in$ *set* (*redex-patterns A*)
unfolding *3 p l-Cs-As* **by** (*metis 3*(*3*) *redex-patterns-elem-rule'*)
then show *?thesis* **by** *simp*
next

```
          case 2
          with i have (γ, p) ∈ set (redex-patterns B)
         unfolding B l-Cs-As l-As-Bs p alpha-beta using redex-patterns-elem-rule'
by blast
          then show ?thesis by simp
        qed
      qed
    }
    moreover
    {fix γ p assume rdp:(γ, p) ∈ set (redex-patterns A) ∪ set (redex-patterns B)
      have (γ, p) ∈ set (redex-patterns C) proof(cases p = [])
        case True
        from rdp lhs have γ = α
          unfolding 3 B alpha-beta True
       by (metis 3(1) B Un-iff alpha-beta list.set-intros(1) option.inject prod.inject
redex-patterns.simps(3) redex-patterns-label subterm.prems(1) subterm.prems(2))
        then show ?thesis unfolding C True by simp
      next
        case False
            from rdp consider (γ, p) ∈ set (redex-patterns A) | (γ, p) ∈ set
(redex-patterns B) by force
        then show ?thesis proof(cases)
          case 1
           then obtain p2 i where i:i < length As i < length (var-poss-list (lhs
α))
          and p:p = var-poss-list (lhs α) ! i @p2 and (γ, p2) ∈ set (redex-patterns
(As!i))
            using 3 redex-patterns-elem-rule False by blast
          with IH have (γ, p2) ∈ set (redex-patterns (Cs!i)) by blast
          with i l-Cs-As show ?thesis unfolding C p
            by (metis redex-patterns-elem-rule')
        next
          case 2
           then obtain p2 i where i:i < length Bs i < length (var-poss-list (lhs
α))
          and p:p = var-poss-list (lhs α) ! i @p2 and (γ, p2) ∈ set (redex-patterns
(Bs!i))
            using B alpha-beta redex-patterns-elem-rule False by blast
           with IH have (γ, p2) ∈ set (redex-patterns (Cs!i)) using l-As-Bs by
simp
          with i l-Cs-As l-As-Bs show ?thesis unfolding C p
            by (metis redex-patterns-elem-rule')
        qed
      qed
    }
    ultimately show ?thesis by auto
  qed
 qed
qed
```

lemma *redex-patterns-join-list*:
  assumes *join-list As* = *Some A* **and** ∀ *a* ∈ *set As*. *a* ∈ *wf-pterm R*
  shows *set* (*redex-patterns A*) = ⋃ (*set* (*map* (*set* ∘ *redex-patterns*) *As*))
  using *assms* **proof**(*induct As arbitrary*:*A*)
  case (*Cons a As*)
  show *?case* **proof**(*cases As* = [])
    case *True*
    from *Cons*(*2*,*3*) **have** *a* = *A*
      unfolding *True join-list.simps* **by** (*simp add*: *join-with-source*)
    then show *?thesis* **unfolding** *True* **by** *simp*
  next
    case *False*
    then have ∗:*join-list* (*a#As*) = *join-opt a* (*join-list As*)
      using *join-list.elims* **by** *blast*
    with *Cons*(*2*) **obtain** *A′* **where** *A′*:*join-list As* = *Some A′* **by** *fastforce*
      with *Cons*(*1*,*3*) **have** *set* (*redex-patterns A′*) = ⋃ (*set* (*map* (*set* ∘ *redex-patterns*) *As*))
      by *simp*
   then show *?thesis* **using** *redex-patterns-join* ∗ *Cons*(*2*,*3*) **unfolding** *A′ join-opt.simps*
      by (*metis* (*no-types*, *opaque-lifting*) *A′ Sup-insert insert-iff join-list-wf-pterm list.set*(*2*) *list.simps*(*9*) *o-apply*)
  qed
qed *simp*

lemma *redex-patterns-context*:
  assumes *p* ∈ *poss s*
  shows *redex-patterns* ((*ctxt-of-pos-term p* (*to-pterm s*)) ⟨*A*⟩) = *map* (λ(α, *q*). (α,*p*@*q*)) (*redex-patterns A*)
  using *assms* **proof**(*induct p arbitrary*:*s*)
  case (*Cons i p′*)
  from *Cons*(*2*) **obtain** *f ss* **where** *s*:*s* = *Fun f ss*
    by (*meson args-poss*)
  from *Cons*(*2*) **have** *i*:*i* < *length ss* **and** *p′*:*p′* ∈ *poss* (*ss*!*i*)
    unfolding *s* **by** *auto*
  with *Cons*(*1*) **have** *IH*:*redex-patterns* (*ctxt-of-pos-term p′* (*to-pterm* (*ss*!*i*)))⟨*A*⟩ =
    map (λ(α, *q*). (α,*p′*@*q*)) (*redex-patterns A*) **by** *simp*
   from *i* **have** *l*:*length* (*take i* (*map to-pterm ss*) @ (*ctxt-of-pos-term p′* (*map to-pterm ss* ! *i*))⟨*A*⟩ # *drop* (*Suc i*) (*map to-pterm ss*)) = *length ss*
    by *simp*
  let *?take-i*=*take i* (*map to-pterm ss*)
  let *?ith*=(*ctxt-of-pos-term p′* (*map to-pterm ss* ! *i*))⟨*A*⟩
  let *?drop-i*=*drop* (*Suc i*) (*map to-pterm ss*)
  let *?xs*=*take i* (*map to-pterm ss*) @ (*ctxt-of-pos-term p′* (*map to-pterm ss* ! *i*))⟨*A*⟩ # *drop* (*Suc i*) (*map to-pterm ss*)
  let *?zip*=*zip* [*0*..<*length ss*] (*map redex-patterns ?xs*)
  from *i* **have** *l-zip*:*length ?zip* = *length ss* **by** *auto*
  let *?zip1*=*zip* [*0*..<*i*] (*map redex-patterns ?take-i*)

let *?zip2=zip* [*Suc i..<length ss*] (*map redex-patterns ?drop-i*)
have *zip*:*?zip* = *?zip1* @ ((*i*, *redex-patterns ?ith*) # *?zip2*)
unfolding *map-append zip-append2* **using** *i* **by** (*simp add*: *upt-conv-Cons*)
{**fix** *j* **assume** *j*:*j* < *length* (*map* (λ(*x*, *y*). *map* (λ(α, *p*). (α, *x* # *p*)) *y*) *?zip1*)
with *i* **have** (*map redex-patterns ?take-i*)!*j* = []
by (*simp add*: *redex-patterns-to-pterm*)
with *j* **have** *?zip1* ! *j* = (*j*, [])
by *simp*
with *j* **have** *map* (λ(*x*, *y*). *map* (λ(α, *p*). (α, *x* # *p*)) *y*) *?zip1* ! *j* = []
by *simp*
}
then have *1*:*concat* (*map* (λ(*x*, *y*). *map* (λ(α, *p*). (α, *x* # *p*)) *y*) *?zip1*) = []
by (*metis length-0-conv length-greater-0-conv less-nat-zero-code nth-concat-split*)

{**fix** *j* **assume** *j*:*j* < *length* (*map* (λ(*x*, *y*). *map* (λ(α, *p*). (α, *x* # *p*)) *y*) *?zip2*)
with *i* **have** (*map redex-patterns ?drop-i*)!*j* = []
by (*simp add*: *redex-patterns-to-pterm*)
with *j* **have** *?zip2* ! *j* = (*j*+*Suc i*, [])
by *simp*
with *j* **have** *map* (λ(*x*, *y*). *map* (λ(α, *p*). (α, *x* # *p*)) *y*) *?zip2* ! *j* = []
by *simp*
}
then have *2*:*concat* (*map* (λ(*x*, *y*). *map* (λ(α, *p*). (α, *x* # *p*)) *y*) *?zip2*) = []
by (*metis length-0-conv length-greater-0-conv less-nat-zero-code nth-concat-split*)

show *?case* **unfolding** *s to-pterm.simps ctxt-of-pos-term.simps intp-actxt.simps redex-patterns.simps l zip*
unfolding *map-append concat-append 1 list.map*(*2*) *concat.simps 2* **using** *IH i* **by** *simp*
qed *simp*

lemma *redex-patterns-prule*:
assumes *l*:*length ts* = *length* (*var-poss-list* (*lhs* α))
shows *redex-patterns* (*Prule* α (*map to-pterm ts*)) = [(α, [])]
proof−
{**fix** *x* **assume** *x*:*x* ∈ *set* (*map2* (λ*x*. *map* (λ(α, *p2*). (α, *x* @ *p2*))) (*var-poss-list* (*lhs* α)) (*map redex-patterns* (*map to-pterm ts*)))
from *l* **have** *length* (*map2* (λ*x*. *map* (λ(α, *p2*). (α, *x* @ *p2*))) (*var-poss-list* (*lhs* α)) (*map redex-patterns* (*map to-pterm ts*))) = *length* (*var-poss-list* (*lhs* α))
by *simp*
with *x* **obtain** *i* **where** *i*:*i* < *length* (*var-poss-list* (*lhs* α)) *x* = (*map2* (λ*x*. *map* (λ(α, *p2*). (α, *x* @ *p2*))) (*var-poss-list* (*lhs* α)) (*map redex-patterns* (*map to-pterm ts*)))!*i*
by (*metis in-set-idx*)
from *i l* **have** *x* = []
using *redex-patterns-to-pterm* **by** *simp*
}
then show *?thesis*
unfolding *redex-patterns.simps* **using** *concat-eq-Nil-conv* **by** *blast*

qed

lemma *redex-patterns-single*:
assumes $p \in$ *poss s* **and** *to-rule* $\alpha \in R$
shows *redex-patterns* (*ll-single-redex s p* α) = [(α, *p*)]
proof−
let *?As*=*map* (*to-pterm* ∘ (λ*pi. s* |- (*p* @ *pi*))) (*var-poss-list* (*lhs* α))
let *?A*=*Prule* α *?As*
have *redex-patterns ?A* = [(α, [])]
using *redex-patterns-prule* **using** *length-map* **by** *fastforce*
moreover have *set* (*redex-patterns* (*ll-single-redex s p* α)) = *set* (*map* (λ (α, *q*). (α,*p*@*q*)) (*redex-patterns ?A*))
using *redex-patterns-context assms redex-patterns-to-pterm*[*of s*] **unfolding** *ll-single-redex-def* **using** *p-in-poss-to-pterm* **by** *fastforce*
ultimately have *set*:*set* (*redex-patterns* (*ll-single-redex s p* α)) = {(α, *p*)}
by *simp*
have *wf*:*ll-single-redex s p* $\alpha \in$ *wf-pterm R*
using *assms left-lin left-linear-trs-def single-redex-wf-pterm* **by** *fastforce*
have *sorted*:*sorted-wrt* (*ord.lexordp* (<)) (*map snd* [(α, *p*)]) **by** *simp*
show *?thesis* **using** *redex-patterns-equal*[*OF wf sorted*] *set* **by** *simp*
qed

lemma *get-label-imp-rdp*:
assumes *get-label* (*labeled-source A* |- *p*) = *Some* (α, *0*)
and $A \in$ *wf-pterm R*
and $p \in$ *poss* (*labeled-source A*)
shows (α, *p*) $\in$ *set* (*redex-patterns A*)
using *assms* **proof**(*induct A arbitrary*:*p*)
case (*Pfun f As*)
then show *?case* **proof**(*cases p*)
case (*Cons i p'*)
from *Pfun*(*4*) **have** *i*:*i* < *length As*
unfolding *Cons* **by** *simp*
moreover from *Pfun*(*2*,*4*) **have** *get-label* (*labeled-source* (*As*!*i*) |- *p'*) = *Some* (α, *0*)
unfolding *Cons* **by** *simp*
moreover from *Pfun*(*4*) **have** $p' \in$ *poss* (*labeled-source* (*As*!*i*))
unfolding *Cons* **using** *i* **by** *simp*
ultimately have (α, *p'*) $\in$ *set* (*redex-patterns* (*As*!*i*))
using *Pfun*(*1*,*3*) **using** *nth-mem* **by** *blast*
then show *?thesis*
unfolding *Cons redex-patterns.simps* **using** *i* **by** (*metis redex-patterns.simps*(*2*) *redex-patterns-elem-fun'*)
qed *simp*
next
case (*Prule* β *As*)
from *Prule*(*3*) **obtain** *f ts* **where** *lhs*:*lhs* β = *Fun f ts*
by (*metis Inl-inject Term.term.simps*(*4*) *case-prodD is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *no-var-lhs term.collapse*(*2*) *term.sel*(*2*) *wf-pterm.simps*)

```
  then show ?case proof(cases p)
    case Nil
    from Prule(2,3,4) show ?thesis
     unfolding Nil labeled-source.simps lhs label-term.simps eval-term.simps subt-at.simps
get-label.simps by simp
  next
    case (Cons i′ p′)
    from Prule(3) have l:length As = length (var-poss-list (lhs β))
       by (metis Inl-inject is-Prule.simps(1) is-Prule.simps(3) length-var-poss-list
length-var-rule term.distinct(1) term.inject(2) wf-pterm.simps)
     from Prule obtain i p2 where p:p = var-poss-list (lhs β)!i @ p2 and i:i <
length As and p2:p2 ∈ poss (labeled-source (As!i))
     by (smt (verit) labeled-source-to-term left-lin-no-var-lhs.get-label-Prule left-lin-no-var-lhs-axioms
list.distinct(1) local.Cons poss-term-lab-to-term)
    let ?x=vars-term-list (lhs β) ! i
    let ?p1=var-poss-list (lhs β) ! i
    have p1:?p1 ∈ poss (labeled-lhs β)
     by (metis i l label-term-to-term nth-mem poss-term-lab-to-term var-poss-imp-poss
var-poss-list-sound)
    have labeled-lhs β |- ?p1 = Var ?x
     using i l by (metis length-var-poss-list var-poss-list-labeled-lhs vars-term-list-labeled-lhs
vars-term-list-var-poss-list)
    then have labeled-source (Prule β As) |- ?p1 = labeled-source (As!i)
      unfolding labeled-source.simps subt-at-subst[OF p1]
     by (smt (verit) Inl-inject Prule.prems(2) apply-lhs-subst-var-rule comp-eq-dest-lhs
eval-term.simps(1) i is-Prule.simps(1) is-Prule.simps(3)
          l length-map length-remdups-eq length-rev length-var-poss-list map-nth-conv
rev-rev-ident term.distinct(1) term.inject(2) wf-pterm.simps)
    with Prule(2,4) have get-label (labeled-source (As!i)|-p2) = Some (α, 0)
      unfolding p labeled-source.simps by auto
    then have (α, p2) ∈ set (redex-patterns (As!i))
      using Prule(1)[of As!i p2] p2 Prule(3) i by (meson fun-well-arg nth-mem)
    then show ?thesis unfolding p redex-patterns.simps using i
      by (metis l redex-patterns.simps(3) redex-patterns-elem-rule′)
  qed
qed simp

lemma redex-pattern-proof-term-equality:
  assumes A ∈ wf-pterm R B ∈ wf-pterm R
    and set (redex-patterns A) = set (redex-patterns B)
    and source A = source B
  shows A = B
  using assms proof(induct A arbitrary:B)
  case (1 x)
  then show ?case
    using redex-poss-empty-imp-empty-step source-empty-step by force
next
  case (2 As f)
  then show ?case proof(cases B)
```

case (*Pfun g Bs*)
from *2*(*4*) **have** *f*:*f* = *g*
unfolding *Pfun* **by** *fastforce*
from *2*(*4*) **have** *len*:*length As* = *length Bs*
unfolding *Pfun f* **by** (*metis length-map source.simps*(*2*) *term.inject*(*2*))
{**fix** *i* **assume** *i*:*i* < *length As*
have *set* (*redex-patterns* (*As*!*i*)) = *set* (*redex-patterns* (*Bs*!*i*)) **proof**(*rule ccontr*)
assume *set* (*redex-patterns* (*As*!*i*)) ≠ *set* (*redex-patterns* (*Bs*!*i*))
then consider ∃ *r*. *r* ∈ *set* (*redex-patterns* (*As*!*i*)) ∧ *r* ∉ *set* (*redex-patterns* (*Bs*!*i*)) |
∃ *r*. *r* ∈ *set* (*redex-patterns* (*Bs*!*i*)) ∧ *r* ∉ *set* (*redex-patterns* (*As*!*i*))
by *blast*
then show *False* **proof**(*cases*)
case *1*
then obtain *p* α **where** (α, *p*) ∈ *set* (*redex-patterns* (*As*!*i*)) **and** *B*:(α, *p*) ∉ *set* (*redex-patterns* (*Bs*!*i*))
by *force*
then have (α, *i*#*p*) ∈ *set* (*redex-patterns* (*Pfun f As*))
by (*meson i redex-patterns-elem-fun′*)
moreover from *B* **have** (α, *i*#*p*) ∉ *set* (*redex-patterns* (*Pfun f Bs*))
by (*metis list.inject redex-patterns-elem-fun*)
ultimately show *?thesis*
using *2.prems*(*2*) *Pfun f* **by** *blast*
next
case *2*
then obtain *p* α **where** (α, *p*) ∈ *set* (*redex-patterns* (*Bs*!*i*)) **and** *A*:(α, *p*) ∉ *set* (*redex-patterns* (*As*!*i*))
by *force*
then have (α, *i*#*p*) ∈ *set* (*redex-patterns* (*Pfun f Bs*))
by (*metis i len redex-patterns-elem-fun′*)
moreover from *A* **have** (α, *i*#*p*) ∉ *set* (*redex-patterns* (*Pfun f As*))
by (*metis list.inject redex-patterns-elem-fun*)
ultimately show *?thesis*
using *2.prems*(*2*) *Pfun f* **by** *blast*
qed
qed
moreover have (*Bs*!*i*) ∈ *wf-pterm R*
using *2*(*2*) *Pfun i len* **by** *auto*
ultimately have *As*!*i* = *Bs*!*i*
using *2*(*1*,*4*) **by** (*metis Pfun i len nth-map nth-mem source.simps*(*2*) *term.inject*(*2*))
}
then show *?thesis*
unfolding *Pfun f* **using** *len* **using** *nth-equalityI* **by** *blast*
next
case (*Prule* α *Bs*)
with *2*(*3*) **show** *?thesis*

by (*metis list.distinct*(*1*) *list.set-intros*(*1*) *redex-patterns.simps*(*3*) *redex-patterns-elem-fun*)

qed *simp*
next
case (*3* α *As*)
then show *?case* **proof**(*cases B*)
case (*Pfun g Bs*)
with *3*(*5*) **show** *?thesis*
by (*metis list.distinct*(*1*) *list.set-intros*(*1*) *redex-patterns.simps*(*3*) *redex-patterns-elem-fun*)
next
case (*Prule* β *Bs*)
from *3*(*5*) **have** α:$\alpha = \beta$
unfolding *Prule* **using** *distinct-snd-rdp*
by (*metis 3.prems*(*1*) *Pair-inject Prule left-lin-no-var-lhs.redex-patterns-label left-lin-no-var-lhs-axioms list.set-intros*(*1*) *option.inject redex-patterns.simps*(*3*))
from *3* **have** *len*:*length As* = *length Bs*
using *Prule* α **by** (*metis length-args-well-Prule wf-pterm.intros*(*3*))
have *len2*:*length* (*var-poss-list* (*lhs* β)) = *length Bs*
by (*metis 3.hyps*(*1*) *3.hyps*(*2*) α *len length-var-poss-list length-var-rule*)
{**fix** *i* **assume** *i*:$i <$ *length As*
obtain *pi* **where** *pi*:*var-poss-list* (*lhs* β) ! i = *pi*
by *auto*
have *set* (*redex-patterns* (*As*!i)) = *set* (*redex-patterns* (*Bs*!i)) **proof**(*rule ccontr*)
assume *set* (*redex-patterns* (*As*!i)) $\neq$ *set* (*redex-patterns* (*Bs*!i))
then consider $\exists\, r.\ r \in$ *set* (*redex-patterns* (*As*!i)) $\wedge\ r \notin$ *set* (*redex-patterns* (*Bs*!i)) |
$\exists\, r.\ r \in$ *set* (*redex-patterns* (*Bs*!i)) $\wedge\ r \notin$ *set* (*redex-patterns* (*As*!i))
by *blast*
then show *False* **proof**(*cases*)
case *1*
then obtain *p* β **where** (β, *p*) $\in$ *set* (*redex-patterns* (*As*!i)) **and** *B*:(β, *p*) $\notin$ *set* (*redex-patterns* (*Bs*!i))
by *force*
then show *False*
using *3*(*4*,*5*) **by** (*metis Prule* α *i len len2 redex-patterns-elem-rule*$'$ *redex-patterns-rule*$''$)
next
case *2*
then obtain *p* β **where** (β, *p*) $\in$ *set* (*redex-patterns* (*Bs*!i)) **and** *A*:(β, *p*) $\notin$ *set* (*redex-patterns* (*As*!i))
by *force*
then show *False*
using *3* **by** (*metis Prule* α *i len len2 redex-patterns-elem-rule*$'$ *redex-patterns-rule*$''$ *wf-pterm.intros*(*3*))
qed
qed
moreover have (*Bs*!i) $\in$ *wf-pterm R*

using *3.prems*(*1*) *Prule i len* **by** *auto*
moreover have *co-initial* (*As*!*i*) (*Bs*!*i*)
using *3* **by** (*metis Prule α co-init-prule i wf-pterm.intros*(*3*))
ultimately have *As*!*i* = *Bs*!*i*
using *3*(*3*) **by** (*simp add*: *i*)
}
then show *?thesis*
unfolding *Prule α* **using** *len* **using** *nth-equalityI* **by** *blast*
qed *simp*
qed

end

abbreviation *single-steps* :: (*'f*, *'v*) *pterm* ⇒ (*'f*, *'v*) *pterm list*
where *single-steps A* ≡ *map* (λ (α, *p*). *ll-single-redex* (*source A*) *p* α) (*redex-patterns A*)

context *left-lin-wf-trs*
begin

lemma *ll-no-var-lhs*: *left-lin-no-var-lhs R*
by (*simp add*: *left-lin-axioms left-lin-no-var-lhs-def no-var-lhs-axioms*)

lemma *single-step-redex-patterns*:
assumes *A* ∈ *wf-pterm R* Δ ∈ *set* (*single-steps A*)
shows ∃ *p* α. Δ = *ll-single-redex* (*source A*) *p* α ∧ (α, *p*) ∈ *set* (*redex-patterns A*) ∧ *redex-patterns* Δ = [(α, *p*)]
proof−
from *assms* **obtain** *p* α **where** Δ:Δ = *ll-single-redex* (*source A*) *p* α **and** *rdp*:(α, *p*) ∈ *set* (*redex-patterns A*)
by *auto*
moreover have *to-rule* α ∈ *R*
using *rdp assms*(*1*) *labeled-wf-pterm-rule-in-TRS left-lin-no-var-lhs.redex-patterns-label ll-no-var-lhs* **by** *fastforce*
moreover have *p* ∈ *poss* (*source A*)
using *assms rdp left-lin-no-var-lhs.redex-patterns-label ll-no-var-lhs* **by** *blast*
ultimately show *?thesis*
using Δ *left-lin-no-var-lhs.redex-patterns-single*[*OF ll-no-var-lhs*] **by** *blast*
qed

lemma *single-step-wf*:
assumes *A* ∈ *wf-pterm R* **and** Δ ∈ *set* (*single-steps A*)
shows Δ ∈ *wf-pterm R*
proof−
from *assms* **obtain** *p* α **where** *p*:*p* ∈ *poss* (*source A*) Δ = *ll-single-redex* (*source A*) *p* α **and** *get-label* ((*labeled-source A*)|-*p*) = *Some* (α, *0*)
using *left-lin-no-var-lhs.redex-patterns-label left-lin-no-var-lhs.redex-patterns-subset-possL possL-subset-poss-source ll-no-var-lhs* **by** *fastforce*

then have *to-rule* $\alpha \in R$
using *assms*(*1*) *labeled-wf-pterm-rule-in-TRS* **by** *fastforce*
with *p* **show** *?thesis* **using** *single-redex-wf-pterm*
using *left-lin left-linear-trs-def* **by** *fastforce*
qed

lemma *source-single-step*:
assumes Δ:$\Delta \in$ *set* (*single-steps A*) **and** *wf*:*A* $\in$ *wf-pterm R*
shows *source* Δ = *source A*
proof−
let *?s*=*source A*
from Δ **obtain** *p* α **where** *pa*:Δ = *ll-single-redex ?s p* α (α, *p*) $\in$ *set* (*redex-patterns A*)
by *auto*
from *pa* **have** *lab-p*:*get-label* (*labeled-source A* |-*p*) = *Some* (α, *0*) **and** *p*:*p* $\in$ *poss ?s*
using *left-lin-no-var-lhs.redex-patterns-label ll-no-var-lhs wf* **by** *blast+*
from *lab-p p* **obtain** *p′* **where** *p′*:*p′*$\in$*poss A* **and** *ctxt*:*ctxt-of-pos-term p* (*source A*) = *source-ctxt* (*ctxt-of-pos-term p′ A*)
and *Ap′*:*A* |- *p′* = *Prule* α (*map* (λ*i*. *A* |- (*p′* @ [*i*])) [*0*..<*length* (*var-rule* α)])
using *poss-labeled-source wf* **by** *force*
have *l*:*length* (*var-rule* α) = *length* (*var-poss-list* (*lhs* α))
using *wf* **by** (*metis Ap′ Inl-inject Term.term.simps*(*4*) *is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *length-var-poss-list length-var-rule p′ subt-at-is-wf-pterm term.inject*(*2*) *wf-pterm.simps*)
{**fix** *i* **assume** *i*:*i* < *length* (*var-rule* α)
let *?pi*=*var-poss-list* (*lhs* α)!*i*
obtain *x* **where** *x*:*lhs* α |- *?pi* = *Var x var-rule* α!*i* = *x*
by (*metis comp-apply i l length-remdups-eq length-rev length-var-poss-list rev-rev-ident vars-term-list-var-poss-list*)
have *?s*|-*p* = *lhs* α · ⟨*map source* (*map* (λ*i*. *A* |- (*p′* @ [*i*])) [*0*..<*length* (*var-rule* α)])⟩$_\alpha$
using *Ap′ ctxt* **by** (*metis ctxt-of-pos-term-well ctxt-supt-id local.wf p p′ replace-at-subt-at source.simps*(*3*) *source-ctxt-apply-term*)
moreover have *lhs* α · ⟨*map source* (*map* (λ*i*. *A* |- (*p′* @ [*i*])) [*0*..<*length* (*var-rule* α)])⟩$_\alpha$ |- *?pi* = *map source* (*map* (λ*i*. *A* |- (*p′* @ [*i*])) [*0*..<*length* (*var-rule* α)])!*i*
using *x* **by** (*smt* (*verit*, *ccfv-SIG*) *diff-zero eval-term.simps*(*1*) *i l length-upt lhs-subst-var-i map-eq-imp-length-eq map-nth nth-mem subt-at-subst var-poss-imp-poss var-poss-list-sound*)
ultimately have *?s*|-(*p*@*?pi*) = *source* (*A* |- (*p′* @ [*i*])) **using** *i p* **by** *auto*
then have *map source* (*map* (λ*i*. *A* |- (*p′* @ [*i*])) [*0*..<*length* (*var-rule* α)]) !*i* = *map* (λ*pi*. *source A* |- (*p* @ *pi*)) (*var-poss-list* (*lhs* α)) ! *i*
using *l i* **by** *auto*
}
with *l* **have** *map source* (*map* (λ*i*. *A* |- (*p′* @ [*i*])) [*0*..<*length* (*var-rule* α)]) = *map* (λ*pi*. *source A* |- (*p* @ *pi*)) (*var-poss-list* (*lhs* α))
by (*simp add*: *map-equality-iff*)
then have *source* (*A*|-*p′*) = *lhs* α · ⟨*map* (λ*pi*. *?s* |- (*p* @ *pi*)) (*var-poss-list* (*lhs* α))⟩$_\alpha$

unfolding *Ap′ source.simps* **by** *simp*
with *ctxt* **show** *?thesis* **unfolding** *pa*(*1*) *source-single-redex*[*OF p*] **using** *p′*
by (*metis ctxt-of-pos-term-well ctxt-supt-id wf source-ctxt-apply-term*)
qed

lemma *single-redex-single-step*:
assumes Δ:Δ = *ll-single-redex s p α*
and *p*:*p* ∈ *poss s* **and** *α*:*to-rule α* ∈ *R*
and *src*:*source* Δ = *s*
shows *single-steps* Δ = [Δ]
using *src* **unfolding** Δ *left-lin-no-var-lhs.redex-patterns-single*[*OF ll-no-var-lhs p α*] **by** *simp*

lemma *single-step-label-imp-label*:
assumes Δ:Δ ∈ *set* (*single-steps A*) **and** *q*:*q* ∈ *poss* (*labeled-source* Δ) **and** *wf*:*A* ∈ *wf-pterm R*
and *lab*:*get-label* (*labeled-source* Δ|-*q*) = *Some l*
shows *get-label* (*labeled-source A* |-*q*) = *Some l*
proof−
let *?s*=*source A*
from Δ **obtain** *p α* **where** *pa*:Δ = *ll-single-redex ?s p α* (*α*, *p*) ∈ *set* (*redex-patterns A*)
by *auto*
from *pa* **have** *lab-p*:*get-label* (*labeled-source A* |-*p*) = *Some* (*α*, *0*) **and** *p*:*p* ∈ *poss* (*source A*)
using *left-lin-no-var-lhs.redex-patterns-label ll-no-var-lhs wf* **by** *blast*+
from *pa lab* **obtain** *q′* **where** *l*:*l* = (*α*, *size q′*) **and** *q′*:*q* = *p*@*q′ q′* ∈ *fun-poss* (*lhs α*)
using *single-redex-label*[*OF pa*(*1*) *p*] *q pa*(*2*) *wf*
by (*metis labeled-source-to-term labeled-wf-pterm-rule-in-TRS left-lin-no-var-lhs.redex-patterns-label ll-no-var-lhs poss-term-lab-to-term prod.collapse*)
from *lab-p p* **obtain** *p′* **where** *p′*∈*poss A* **and** *ctxt-of-pos-term p* (*source A*) = *source-ctxt* (*ctxt-of-pos-term p′ A*) **and** *A* |- *p′* = *Prule α* (*map* (*λi. A* |- (*p′* @ [*i*])) [*0*..<*length* (*var-rule α*)])
using *poss-labeled-source wf* **by** *force*
then have *labeled-source A* = (*ctxt-of-pos-term p* (*labeled-source A*))⟨*labeled-source* (*Prule α* (*map* (*λi. A* |- (*p′* @ [*i*])) [*0*..<*length* (*var-rule α*)])))⟩
using *label-source-ctxt p wf* **by** (*metis ctxt-supt-id*)
then have *labeled-source A*|-*q* = *labeled-lhs α* · ⟨*map labeled-source* (*map* (*λi. A* |- (*p′* @ [*i*])) [*0*..<*length* (*var-rule α*)])⟩$_\alpha$ |-*q′*
unfolding *q′ labeled-source.simps* **by** (*metis labeled-source.simps*(*3*) *labeled-source-to-term p poss-term-lab-to-term subt-at-append subt-at-ctxt-of-pos-term*)
then have *get-label* (*labeled-source A*|-*q*) = *Some* (*α*, *size q′*)
using *q′*(*2*) **by** (*simp add*: *label-term-increase*)
with *l* **show** *?thesis* **by** *simp*
qed

lemma *single-steps-measure*:
assumes Δ*1* :Δ*1* ∈ *set* (*single-steps A*) **and** Δ*2*:Δ*2* ∈ *set* (*single-steps A*)

and *wf*:$A \in$ *wf-pterm R* **and** *neq*:$\Delta 1 \neq \Delta 2$
shows *measure-ov* $\Delta 1$ $\Delta 2$ = *0*
proof−
let *?s=source A*
from $\Delta 1$ **obtain** *p* α **where** *pa*:$\Delta 1$ = *ll-single-redex ?s p* α (α, *p*) $\in$ *set* (*redex-patterns A*)
by *auto*
from $\Delta 2$ **obtain** *q* β **where** *qb*:$\Delta 2$ = *ll-single-redex ?s q* β (β, *q*) $\in$ *set* (*redex-patterns A*)
by *auto*
from *neq* **have** *pq*:$p \neq q \vee \alpha \neq \beta$
using *pa*(*1*) *qb*(*1*) **by** *force*
{**assume** *measure-ov* $\Delta 1$ $\Delta 2 \neq$ *0*
then obtain *r* **where** *r1*:$r \in$ *possL* $\Delta 1$ **and** *r2*:$r \in$ *possL* $\Delta 2$
by (*metis card.empty disjoint-iff*)
from *r1* **obtain** p' **where** p':$r = p@p'$ **and** *l1*:*get-label* (*labeled-source* $\Delta 1$ |-*r*) = *Some* (α, *size* p')
using *single-redex-label*[*OF pa*(*1*)] *wf*
by (*smt* (*verit*, *ccfv-SIG*) *labeled-source-to-term labeled-wf-pterm-rule-in-TRS left-lin-no-var-lhs.redex-patterns-label ll-no-var-lhs pa*(*2*) *possL-obtain-label possL-subset-poss-source poss-term-lab-to-term subsetD*)
from *r2* **obtain** q' **where** q':$r = q@q'$ **and** *l2*:*get-label* (*labeled-source* $\Delta 2$ |-*r*) = *Some* (β, *size* q')
using *single-redex-label*[*OF qb*(*1*)] *wf*
by (*smt* (*verit*, *ccfv-SIG*) *labeled-source-to-term labeled-wf-pterm-rule-in-TRS left-lin-no-var-lhs.redex-patterns-label ll-no-var-lhs qb*(*2*) *possL-obtain-label possL-subset-poss-source poss-term-lab-to-term subsetD*)
from *l1* **have** *get-label* (*labeled-source A* |-*r*) = *Some* (α, *size* p')
using $\Delta 1$ *labelposs-subs-poss wf r1 single-step-label-imp-label* **by** *blast*
moreover from *l2* **have** *get-label* (*labeled-source A* |-*r*) = *Some* (β, *size* q')
using $\Delta 2$ *labelposs-subs-poss wf r2 single-step-label-imp-label* **by** *blast*
moreover from *pq* **have** $p' \neq q' \vee \alpha \neq \beta$
using p' q' **by** *blast*
ultimately have *False* **using** p' q' **by** *auto*
}
then show *?thesis* **by** *auto*
qed

lemma *single-steps-orth*:
assumes $\Delta 1$:$\Delta 1 \in$ *set* (*single-steps A*) **and** $\Delta 2$:$\Delta 2 \in$ *set* (*single-steps A*) **and** *wf*:$A \in$ *wf-pterm R*
shows $\Delta 1 \perp_p \Delta 2$
using *single-steps-measure*[*OF* $\Delta 1$ $\Delta 2$ *wf*] *equal-imp-orthogonal*
by (*metis* $\Delta 1$ $\Delta 2$ *ll-no-var-lhs local.wf measure-zero-imp-orthogonal single-step-wf source-single-step*)

lemma *redex-patterns-below*:
assumes *wf*:$A \in$ *wf-pterm R*
and (α, *p*) $\in$ *set* (*redex-patterns A*)

and $(\beta,\ p@q) \in$ *set* (*redex-patterns* A) **and** $q \neq []$
shows $q \notin$ *fun-poss* (*lhs* α)
proof−
let *?Δ1=ll-single-redex* (*source* A) p α
let *?Δ2=ll-single-redex* (*source* A) $(p@q)$ β
from *assms* **have** Δ*1*:*?Δ1* $\in$ *set* (*single-steps* A)
by *force*
from *assms* **have** Δ*2*:*?Δ2* $\in$ *set* (*single-steps* A)
by *force*
from *assms*(*1*,*2*) **have** *possL1*:*possL ?Δ1* $= \{p@p' \mid p'.\ p' \in$ *fun-poss* (*lhs* α)$\}$
by (*metis* (*no-types*, *lifting*) *left-lin-no-var-lhs.redex-pattern-rule-symbol left-lin-no-var-lhs.redex-patterns-label ll-no-var-lhs single-redex-possL*)
from *assms*(*1*,*3*) **have** *possL2*:*possL ?Δ2* $= \{(p@q)@p' \mid p'.\ p' \in$ *fun-poss* (*lhs* β)$\}$
using *left-lin.single-redex-possL left-lin-axioms left-lin-no-var-lhs.redex-pattern-rule-symbol left-lin-no-var-lhs.redex-patterns-label ll-no-var-lhs* **by** *blast*
from *assms* **have** *neq*:*?Δ1* $\neq$ *?Δ2*
by (*metis Pair-inject left-lin-no-var-lhs.redex-patterns-label ll-no-var-lhs self-append-conv single-redex-neq*)
from *single-steps-measure*[*OF* Δ*1* Δ*2 wf neq*] **have** *possL ?Δ1* $\cap$ *possL ?Δ2* $= \{\}$
by (*simp add*: *finite-possL*)
moreover have $[] \in$ *fun-poss* (*lhs* β) **proof**−
have *to-rule* $\beta \in R$
using *assms*(*1*) *assms*(*3*) *left-lin-no-var-lhs.redex-pattern-rule-symbol ll-no-var-lhs* **by** *blast*
then show *?thesis*
using *wf-trs-alt wf-trs-imp-lhs-Fun* **by** *fastforce*
qed
ultimately show *?thesis*
unfolding *possL1 possL2* **by** *auto*
qed

lemma *single-steps-singleton*:
assumes *A-wf*:$A \in$ *wf-pterm R* **and** Δ:*single-steps* $A = [\Delta]$
shows $A = \Delta$
proof−
obtain p α **where** *rdp-*Δ:Δ = *ll-single-redex* (*source* A) p α $(\alpha,\ p) \in$ *set* (*redex-patterns* A) *redex-patterns* $\Delta = [(\alpha,\ p)]$
using *single-step-redex-patterns*[*OF A-wf*] Δ **by** *auto*
then have *rdp-A*:*redex-patterns* $A = [(\alpha,\ p)]$
by (*smt* (*verit*) Δ *in-set-simps*(*2*) *list.map-disc-iff map-eq-Cons-D*)
then show *?thesis*
using *left-lin-no-var-lhs.redex-pattern-proof-term-equality*[*OF ll-no-var-lhs A-wf*]
by (*metis A-wf* Δ *list.set-intros*(*1*) *rdp-*Δ(*3*) *single-step-wf source-single-step*)
qed
end

context *left-lin-no-var-lhs*

begin
lemma *measure-ov-imp-single-step-ov*:
assumes *measure-ov A B* $\neq$ *0* **and** *wf*:*A* $\in$ *wf-pterm R*
shows $\exists\,\Delta \in$ *set* (*single-steps A*). *measure-ov* Δ *B* $\neq$ *0*
proof−
from *assms* **obtain** *r* **where** *r1*:*r* $\in$ *possL A* **and** *r2*:*r* $\in$ *possL B*
by (*metis card.empty disjoint-iff*)
then obtain α *n* **where** *lab*:*get-label* (*labeled-source A* |- *r*) = *Some* (α, *n*)
using *possL-obtain-label* **by** *blast*
with *wf r1* **obtain** *r1 r2* **where** *r*:*r* = *r1*@*r2* **and** *lab-r1*:*get-label* (*labeled-source A* |-*r1*) = *Some* (α, *0*) **and** *n*:*length r2* = *n*
by (*metis* (*no-types*, *lifting*) *append-take-drop-id diff-diff-cancel label-term-max-value labelposs-subs-poss length-drop obtain-label-root subsetD*)
from *r r1* **have** *r1-pos*:*r1* $\in$ *poss* (*labeled-source A*)
using *labelposs-subs-poss poss-append-poss* **by** *blast*
then obtain *q* **where** *q*:*q*$\in$*poss A* **and** *ctxt*:*ctxt-of-pos-term r1* (*source A*) = *source-ctxt* (*ctxt-of-pos-term q A*)
and *Aq*:*A* |- *q* = *Prule* α (*map* (λ*i*. *A* |- (*q* @ [*i*])) [*0*..<*length* (*var-rule* α)])
using *poss-labeled-source wf lab-r1* **by** *blast*
with *r r1* **have** *r2-pos*:*r2* $\in$ *poss* (*source* (*Prule* α (*map* (λ*i*. *A* |- (*q* @ [*i*])) [*0*..<*length* (*var-rule* α)])))
by (*metis* (*no-types*, *lifting*) *ctxt-supt-id fun-poss-imp-poss label-source-ctxt labeled-source-to-term labelposs-subs-fun-poss-source local.wf poss-term-lab-to-term r1-pos replace-at-subt-at subterm-poss-conv*)
from *Aq q wf* **have** *Prule* α (*map* (λ*i*. *A* |- (*q* @ [*i*])) [*0*..<*length* (*var-rule* α)]) $\in$ *wf-pterm R*
using *subt-at-is-wf-pterm* **by** *auto*
moreover then have *is-Fun* (*lhs* α) **using** *no-var-lhs*
using *wf-pterm.cases* **by** *fastforce*
moreover from *lab ctxt* **have** *get-label* (*labeled-source* (*Prule* α (*map* (λ*i*. *A* |- (*q* @ [*i*])) [*0*..<*length* (*var-rule* α)])) |-*r2*) = *Some* (α, *n*)
by (*metis* (*no-types*, *lifting*) *Aq ctxt-supt-id label-source-ctxt labeled-source-to-term local.wf poss-term-lab-to-term q r r1-pos replace-at-subt-at subt-at-append*)
ultimately have *r2-funposs*:*r2* $\in$ *fun-poss* (*lhs* α)
using *labeled-poss-in-lhs*[*OF r2-pos*] *n* **by** *blast*
let *?*Δ=*ll-single-redex* (*source A*) *r1* α
from *lab-r1 r1-pos* **have** *rdp*:(α, *r1*) $\in$ *set* (*redex-patterns A*)
using *redex-patterns-label wf* **by** *auto*
then have Δ:*?*Δ $\in$ *set* (*single-steps A*) **by** *force*
from *r2* **have** *measure-ov ?*Δ *B* $\neq$ *0*
by (*smt* (*verit*, *ccfv-threshold*) *rdp labeled-sources-imp-measure-not-zero labeled-wf-pterm-rule-in-TRS labelposs-subs-poss wf mem-Collect-eq option.simps*(*3*) *possL-obtain-label r r1-pos r2-funposs redex-patterns-label rel-simps*(*70*) *single-redex-possL subsetD*)
with Δ **show** *?thesis* **by** *blast*
qed
end

context *left-lin-no-var-lhs*

begin
lemma *label-single-step*:
  assumes $p \in poss$ (*labeled-source A*) $A \in$ *wf-pterm R*
    and *get-label* (*labeled-source A* |- *p*) = *Some* (α, *n*)
  shows $\exists$ *Ai*. *Ai* $\in$ *set* (*single-steps A*) $\wedge$ *get-label* (*labeled-source Ai* |- *p*) = *Some* (α, *n*)
proof−
  let *?p1*=*take* (*length p* − *n*) *p*
  let *?p2*=*drop* (*length p* − *n*) *p*
  let *?xs*=*map* (*to-pterm* ∘ (λ*pi*. (*source A*)|-(*p*@*pi*))) (*var-poss-list* (*lhs* α))
  from *assms*(*1*) **have** *p1-pos*:*?p1* $\in$ *poss* (*labeled-source A*)
    by (*metis append-take-drop-id poss-append-poss*)
  have *lab*:*get-label* (*labeled-source A* |- *?p1*) = *Some* (α, *0*)
    using *obtain-label-root*[*OF assms*(*1*) *assms*(*3*) *assms*(*2*)] **by** *simp*
  with *assms* **have** *rdp*:(α, *?p1*) $\in$ *set* (*redex-patterns A*)
    using *redex-patterns-label*[*OF assms*(*2*)] **by** (*metis labeled-source-to-term obtain-label-root poss-term-lab-to-term*)
  then have *ll-single-redex* (*source A*) *?p1* α $\in$ *set* (*single-steps A*) **by** *force*
  then obtain *Ai* **where** *Ai*:*Ai* $\in$ *set* (*single-steps A*) *Ai* = *ll-single-redex* (*source A*) *?p1* α
    by *presburger*
  from *rdp* **obtain** *p′ As* **where** *p′*:*A*|-*p′* = *Prule* α *As p′* $\in$ *poss A ctxt-of-pos-term ?p1* (*source A*) = *source-ctxt* (*ctxt-of-pos-term p′ A*)
    using *poss-labeled-source*[*OF p1-pos*] *lab assms*(*2*) **by** *blast*
  from *p′ assms*(*2*) **have** *A*|-*p′* $\in$ *wf-pterm R*
    using *subt-at-is-wf-pterm* **by** *blast*
  moreover from *p′ assms* **have** *get-label* (*labeled-source* (*A*|-*p′*) |- *?p2*) = *Some* (α, *n*)
    by (*smt* (*verit*, *ccfv-SIG*) *append-take-drop-id ctxt-supt-id label-source-ctxt p1-pos rdp redex-patterns-label replace-at-subt-at subterm-poss-conv*)
  ultimately have *p2-pos*:*?p2* $\in$ *fun-poss* (*lhs* α)
    using *labeled-poss-in-lhs no-var-lhs assms p′*
    by (*smt* (*verit*, *ccfv-threshold*) *append-take-drop-id case-prod-conv ctxt-of-pos-term-well ctxt-supt-id diff-diff-cancel label-term-max-value*
      labeled-source-to-term labeled-wf-pterm-rule-in-TRS length-drop poss-append-poss poss-term-lab-to-term replace-at-subt-at source-ctxt-apply-term)
  then have *l*:*get-label* (*labeled-source* (*Prule* α *?xs*) |- *?p2*) = *Some* (α, *n*)
    using *label-term-increase assms* **by** (*metis* (*no-types*, *lifting*) *add-0 diff-diff-cancel label-term-max-value labeled-source.simps*(*3*) *length-drop*)
  from *p1-pos* **have** *?p1* $\in$ *poss* (*source A*) **by** *simp*
  then have *get-label* (*labeled-source Ai* |- *p*) = *Some* (α, *n*)
    unfolding *Ai*(*2*) **by** (*metis p2-pos append-take-drop-id l label-ctxt-apply-term label-term-increase labeled-source.simps*(*3*) *ll-single-redex-def*)
  with *Ai* **show** *?thesis*
    by *blast*
qed

lemma *proof-term-matches*:
  assumes $A \in$ *wf-pterm R B* $\in$ *wf-pterm R linear-term A*

and $\bigwedge\alpha\ r.\ (\alpha,\ r) \in$ *set* (*redex-patterns A*) = ((α, r) $\in$ *set* (*redex-patterns B*) $\wedge$ $r \in$ *fun-poss* (*source A*))
and *source A* $\cdot$ σ = *source B*
shows $A \cdot$ (*mk-subst Var* (*match-substs A B*)) = B
proof−
{**fix** *p g ts* **assume** $p \in$ *poss A A*|-*p* = *Fun g ts*
with *assms* **have** $\exists$ *Bs. length ts* = *length Bs* $\wedge$ *B*|-*p* = *Fun g Bs*
using *assms* **proof**(*induct A arbitrary*: *B p rule*:*pterm-induct*)
case (*Pfun f As*)
then have $\bigwedge\alpha$. (α, []) $\notin$ *set* (*redex-patterns* (*Pfun f As*))
by (*meson list.distinct*(*1*) *redex-patterns-elem-fun*)
with *Pfun*(*5*) **have** $\neg$ (*is-Prule B*)
by (*metis empty-pos-in-poss is-FunI is-Prule.elims*(*2*) *list.set-intros*(*1*) *poss-is-Fun-fun-poss redex-patterns.simps*(*3*) *source.simps*(*2*) *subt-at.simps*(*1*))
with *Pfun*(*6*) **obtain** *Bs* **where** *B*:*B* = *Pfun f Bs* **and** *l*:*length Bs* = *length As*
by (*smt* (*verit*, *del-insts*) *eval-term.simps*(*2*) *is-Prule.elims*(*3*) *length-map source.simps*(*1*) *source.simps*(*2*) *term.distinct*(*1*) *term.inject*(*2*))
then show *?case* **proof**(*cases p*)
case *Nil*
from *Pfun*(*8*) **show** *?thesis* **unfolding** *Nil B* **using** *l* **by** *simp*
next
case (*Cons i p'*)
from *Pfun*(*7*) **have** *i*:$i <$ *length As* **and** *p'*:$p' \in$ *poss* (*As*!*i*) **and** *a*:*As*!*i* $\in$ *set As*
unfolding *Cons* **by** *simp-all*
from *Pfun*(*8*) **have** *at-p'*:(*As*!*i*)|-*p'* = *Fun g ts*
unfolding *Cons* **by** *simp*
from *Pfun*(*2*) **have** *a-wf*:*As*!*i* $\in$ *wf-pterm R*
using *i nth-mem* **by** *blast*
from *Pfun*(*3*) **have** *b-wf*:*Bs*!*i* $\in$ *wf-pterm R*
unfolding *B* **using** *i l* **by** *auto*
from *Pfun*(*4*) **have** *a-lin*:*linear-term* (*As*!*i*)
using *i* **by** *simp*
{**fix** α *r* **assume** (α, *r*) $\in$ *set* (*redex-patterns* (*As*!*i*))
then have (α, *i*#*r*) $\in$ *set* (*redex-patterns* (*Pfun f As*))
by (*meson i redex-patterns-elem-fun'*)
with *Pfun*(*5*) **have** (α, *r*) $\in$ *set* (*redex-patterns* (*Bs*!*i*)) $\wedge$ *i*#*r* $\in$ *fun-poss* (*source* (*Pfun f As*))
unfolding *B* **by** (*metis list.inject redex-patterns-elem-fun*)
then have (α, *r*) $\in$ *set* (*redex-patterns* (*Bs*!*i*)) $\wedge$ *r* $\in$ *fun-poss* (*source* (*As*!*i*))
using *i* **by** *simp*
} **moreover** {**fix** α *r* **assume** (α, *r*) $\in$ *set* (*redex-patterns* (*Bs*!*i*)) **and** *r*:*r* $\in$ *fun-poss* (*source* (*As*!*i*))
then have (α, *i*#*r*) $\in$ *set* (*redex-patterns B*)
unfolding *B* **using** *i l* **by** (*metis redex-patterns-elem-fun'*)
moreover from *r* **have** *i*#*r* $\in$ *fun-poss* (*source* (*Pfun f As*))
using *i* **unfolding** *source.simps fun-poss.simps length-map* **by** *simp*

ultimately have $(\alpha, r) \in$ *set* (*redex-patterns* (*As*!*i*))
using *Pfun*(*5*) *i* **by** (*metis list.inject redex-patterns-elem-fun*)
}
ultimately have *rdp*:$\bigwedge\alpha$ *r*. ((α, *r*) $\in$ *set* (*redex-patterns* (*As* ! *i*))) = ((α, *r*) $\in$ *set* (*redex-patterns* (*Bs* ! *i*)) $\wedge$ *r* $\in$ *fun-poss* (*source* (*As* ! *i*)))
by *blast*
from *Pfun*(*6*) **have** *sigma*:*source* (*As*!*i*) $\cdot$ σ = *source* (*Bs*!*i*)
unfolding *B source.simps eval-term.simps* **using** *i l* **using** *map-nth-conv* **by** *fastforce*
from *Pfun*(*1*)[*OF a a-wf b-wf a-lin rdp sigma p′ at-p′ a-wf b-wf a-lin rdp sigma*]
obtain *ss* **where** *length ts* = *length ss* **and** *Bs*!*i* |- *p′* = *Fun g ss* **by** *blast*
then show *?thesis* **unfolding** *B Cons* **using** *i l* **by** *simp*
qed
next
case (*Prule* α *As*)
from *Prule*(*5*) **have** (α, []) $\in$ *set* (*redex-patterns B*)
by *auto*
then obtain *Bs* **where** *B*:*B* = *Prule* α *Bs*
by (*smt* (*verit*, *ccfv-threshold*) *Prule.prems*(*2*) *in-set-idx in-set-simps*(*3*) *redex-patterns-elem-fun less-nat-zero-code list.distinct*(*1*)
nat-neq-iff nth-Cons-0 order-pos.dual-order.refl prod.inject redex-patterns.simps(*1*) *redex-patterns.simps*(*3*) *redex-patterns-order wf-pterm.simps*)
with *Prule*(*2*,*3*) **have** *l*:*length As* = *length Bs*
using *length-args-well-Prule* **by** *blast*
show *?case* **proof**(*cases p*)
case *Nil*
from *Prule*(*8*) **show** *?thesis* **unfolding** *Nil B* **using** *l* **by** *simp*
next
case (*Cons i p′*)
from *Prule*(*7*) **have** *i*:*i* < *length As* **and** *p′*:*p′* $\in$ *poss* (*As*!*i*) **and** *a*:*As*!*i* $\in$ *set As*
unfolding *Cons* **by** *simp-all*
from *Prule*(*8*) **have** *at-p′*:(*As*!*i*)|-*p′* = *Fun g ts*
unfolding *Cons* **by** *simp*
from *Prule*(*2*) **have** *a-wf*:*As*!*i* $\in$ *wf-pterm R*
using *i nth-mem* **by** *blast*
from *Prule*(*3*) **have** *b-wf*:*Bs*!*i* $\in$ *wf-pterm R*
unfolding *B* **using** *i l* **by** *auto*
from *Prule*(*4*) **have** *a-lin*:*linear-term* (*As*!*i*)
using *i* **by** *simp*
let *?pi*=*var-poss-list* (*lhs* α) ! *i*
let *?xi*=*vars-term-list* (*lhs* α) ! *i*
have *i′*:*i* < *length* (*var-poss-list* (*lhs* α))
using *i Prule*(*2*) **by** (*metis Inl-inject is-Prule.simps*(*1*) *is-Prule.simps*(*3*) *length-var-poss-list length-var-rule term.distinct*(*1*) *term.inject*(*2*) *wf-pterm.simps*)

have *eval-lhs′*:$\bigwedge\sigma$. *lhs* α $\cdot$ σ |- *?pi* = σ *?xi*
by (*metis eval-term.simps*(*1*) *i′ length-var-poss-list nth-mem subt-at-subst*

var-poss-imp-poss var-poss-list-sound vars-term-list-var-poss-list)
then have *eval-lhs*:$\bigwedge\sigma$ *q*. *lhs* $\alpha \cdot \sigma$ |- (*?pi*@*q*) = σ *?xi* |- *q*
by (*smt* (*verit*) *i′ nth-mem poss-imp-subst-poss subt-at-append var-poss-imp-poss var-poss-list-sound*)
have *i* < *length* (*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* α)) (*map redex-patterns Bs*))
unfolding *length-map length-zip* **using** *i l i′* **by** *simp*
moreover have *zip* (*var-poss-list* (*lhs* α)) (*map redex-patterns Bs*) ! *i* = (*?pi*, *redex-patterns* (*Bs*!*i*))
using *i i′ l* **by** *force*
ultimately have *map-rdp*:(*map2* (λ*p1*. *map* (λ(α, *p2*). (α, *p1* @ *p2*))) (*var-poss-list* (*lhs* α)) (*map redex-patterns Bs*))!*i* = *map* (λ(α, *p2*). (α, *?pi* @ *p2*)) (*redex-patterns* (*Bs*!*i*))
by *simp*
have *l′*:*length* (*var-rule* α) = *length* (*vars-term-list* (*lhs* α))
using *B Prule.prems*(*2*) *left-lin.length-var-rule left-lin-axioms wf-pterm.simps* **by** *fastforce*
{**fix** β *r* **assume** β*r*:(β, *r*) ∈ *set* (*redex-patterns* (*As*!*i*))
from *i′* **have** (β, *?pi*@*r*) ∈ *set* (*redex-patterns* (*Prule* α *As*))
using *redex-patterns-elem-rule′*[*OF* β*r i*] **by** *simp*
with *Prule*(*5*) **have** *1*:(β, *?pi*@*r*) ∈ *set* (*redex-patterns B*) **and** *2*:*?pi*@*r* ∈ *fun-poss* (*source* (*Prule* α *As*))
by *presburger+*
from *1* **have** (β, *r*) ∈ *set* (*redex-patterns* (*Bs*!*i*))
using *redex-patterns-rule′′* **by** (*metis B Prule.prems*(*2*) *i l*)
moreover have *r* ∈ *fun-poss* (*source* (*As*!*i*))
by (*metis* β*r a-wf get-label-imp-labelposs labeled-source-to-term label-poss-subs-fun-poss-source left-lin-no-var-lhs.redex-patterns-label left-lin-no-var-lhs-axioms option.distinct*(*1*) *poss-term-lab-to-term*)
ultimately have (β, *r*) ∈ *set* (*redex-patterns* (*Bs*!*i*)) ∧ *r* ∈ *fun-poss* (*source* (*As*!*i*))
by *simp*
} **moreover** {**fix** β *r* **assume** β*r*:(β, *r*) ∈ *set* (*redex-patterns* (*Bs*!*i*)) **and** *r*:*r* ∈ *fun-poss* (*source* (*As*!*i*))
let *?x*=*var-rule* α ! *i*
from *l′* **have** *x*:*lhs* α |- *?pi* = *Var ?x*
using *i* **by** (*metis comp-apply eval-lhs′ length-remdups-eq length-rev rev-rev-ident subst-apply-term-empty*)
with *r* **have** *r*:*r* ∈ *fun-poss* (*lhs* α |- *?pi* · ⟨*map source As*⟩$_\alpha$)
using *lhs-subst-var-i l′ i i′* **by** (*metis* (*mono-tags*, *lifting*) *eval-term.simps*(*1*) *length-map length-var-poss-list nth-map*)
from β*r* **have** (β, *?pi*@*r*) ∈ *set* (*redex-patterns B*)
unfolding *B* **using** *i l* **using** *redex-patterns-elem-rule′*[*OF* β*r i*[*unfolded l*] *i′*] **by** *simp*
moreover from *r x* **have** *?pi*@*r* ∈ *fun-poss* (*source* (*Prule* α *As*))
using *i* **unfolding** *source.simps fun-poss.simps*
by (*metis* (*no-types*, *lifting*) *eval-lhs eval-lhs′ fun-poss-fun-conv fun-poss-imp-poss i′ is-FunI nth-mem*
pos-append-poss poss-imp-subst-poss poss-is-Fun-fun-poss subt-at-subst

```
var-poss-imp-poss var-poss-list-sound)
        ultimately have (β, ?pi @ r) ∈ set (redex-patterns (Prule α As))
          using Prule(5) by presburger
        then have (β, r) ∈ set (redex-patterns (As!i))
          using i redex-patterns-rule'' Prule.prems(1) by blast
      }
      ultimately have rdp:⋀α r. ((α, r) ∈ set (redex-patterns (As ! i))) = ((α,
r) ∈ set (redex-patterns (Bs ! i)) ∧ r ∈ fun-poss (source (As ! i)))
        by blast
      from Prule(6) have sigma:source (As!i) · σ = source (Bs!i)
        unfolding B source.simps eval-term.simps using i l map-nth-conv
        by (smt (verit, best) B Inl-inject Prule.prems(2) apply-lhs-subst-var-rule
comp-apply eval-lhs' i' is-Prule.simps(1) is-Prule.simps(3) length-map length-remdups-eq
length-rev length-var-rule nth-mem poss-imp-subst-poss rev-swap subt-at-subst term.distinct(1)
term.inject(2) var-poss-imp-poss var-poss-list-sound wf-pterm.simps)
      from Prule(1)[OF a a-wf b-wf a-lin rdp sigma p' at-p' a-wf b-wf a-lin rdp
sigma]
      obtain ss where length ts = length ss and Bs!i |- p' = Fun g ss by blast
      then show ?thesis unfolding B Cons using i l by simp
    qed
  qed simp
 }
 then show ?thesis using fun-poss-eq-imp-matches[OF assms(3)] by simp
qed
end

context left-lin-wf-trs
begin
lemma join-single-steps-wf:
 assumes A ∈ wf-pterm R
 and As = filter f (single-steps A) and As ≠ []
 shows ∃D. join-list As = Some D ∧ D ∈ wf-pterm R
proof−
 {fix a1 a2 assume a1:a1 ∈ set (single-steps A) and a2:a2 ∈ set (single-steps
A)
   with assms(1,2) have a1 ⊥p a2 ∨ a1 = a2
     using single-steps-orth by presburger
   moreover from a1 have a1 ∈ wf-pterm R
     using single-step-wf[OF assms(1)] assms(2) by presburger
   moreover from a2 have a2 ∈ wf-pterm R
     using single-step-wf[OF assms(1)] assms(2) by presburger
   ultimately have a1 ⊔ a2 ≠ None
     using join-same orth-imp-join-defined no-var-lhs by fastforce
 }
 then show ?thesis using left-lin-no-var-lhs.join-list-defined[OF ll-no-var-lhs]
assms(2,3) single-step-wf[OF assms(1)] by simp
qed

lemma single-steps-join-list:
```

assumes *join-list As* = *Some A* **and** $\forall$ *a* $\in$ *set As*. *a* $\in$ *wf-pterm R*
shows *set* (*single-steps A*) = $\bigcup$ (*set* (*map* (*set* $\circ$ *single-steps*) *As*))
proof−
have *rdp*:*set* (*redex-patterns A*) = $\bigcup$ (*set* (*map* (*set* $\circ$ *redex-patterns*) *As*))
using *left-lin-no-var-lhs.redex-patterns-join-list assms ll-no-var-lhs* **by** *blast*
{**fix** *a* **assume** *a* $\in$ *set* (*single-steps A*)
then obtain α *p* **where** *a*:*a* = *ll-single-redex* (*source A*) *p* α **and** (α, *p*) $\in$ *set* (*redex-patterns A*) **by** *auto*
with *rdp* **obtain** *Ai* **where** *Ai*:*Ai* $\in$ *set As* **and** (α, *p*) $\in$ *set* (*redex-patterns Ai*) **by** *auto*
then have *a* $\in$ *set* (*single-steps Ai*)
unfolding *a* **using** *left-lin-no-var-lhs.source-join-list*[*OF ll-no-var-lhs assms*] **by** *force*
with *Ai* **have** *a* $\in$ $\bigcup$ (*set* (*map* (*set* $\circ$ *single-steps*) *As*)) **by** *auto*
} **moreover**
{**fix** *a* **assume** *a* $\in$ $\bigcup$ (*set* (*map* (*set* $\circ$ *single-steps*) *As*))
then obtain *Ai* **where** *Ai*:*Ai* $\in$ *set As a* $\in$ *set* (*single-steps Ai*)
by (*smt* (*verit*, *best*) *UnionE comp-def in-set-idx length-map map-nth-eq-conv nth-mem*)
then obtain α *p* **where** *a*:*a* = *ll-single-redex* (*source Ai*) *p* α **and** (α, *p*) $\in$ *set* (*redex-patterns Ai*) **by** *auto*
with *rdp Ai* **have** (α, *p*) $\in$ *set* (*redex-patterns A*) **by** *auto*
then have *a* $\in$ *set* (*single-steps A*)
unfolding *a* **using** *left-lin-no-var-lhs.source-join-list*[*OF ll-no-var-lhs assms*] *Ai* **by** *force*
}
ultimately show *?thesis* **by** *fastforce*
qed
end

end

References

[1] C. Kirk and A. Middeldorp. Formalizing simultaneous critical pairs for confluence of left-linear rewrite systems. In *Proc. 14th International Conference on Certified Programs and Proofs*, pages 156–170, 2025.

[2] C. Kohl and A. Middeldorp. A formalization of the development closedness criterion for left-linear term rewrite systems. In *Proc. 12th International Conference on Certified Programs and Proofs*, pages 197–210, 2023.

[3] C. Kohl and A. Middeldorp. Formalizing almost development closed critical pairs (short paper). In *Proc. 14th International Joint Conference on Automated Reasoning*, volume 268 of *LIPIcs*, pages 38:1–38:8, 2023.

[4] C. Kohl and A. Middeldorp. Formalizing confluence and commutation criteria using proof terms. In *Proc. 12th International Workshop on Confluence*, pages 49–54, 2023. Available from http://cl-informatik.uibk.ac.at/iwc/2023/proceedings.pdf.

[5] TeReSe, editor. *Term Rewriting Systems*, volume 55 of *Cambridge Tracts in Theoretical Computer Science*. Cambridge University Press, 2003.

[6] V. van Oostrom and R. de Vrijer. Four equivalent equivalences of reductions. In *Proc. 2nd International Workshop on Reduction Strategies in Rewriting and Programming*, volume 70(6) of *Electronic Notes in Theoretical Computer Science*, pages 21–61, 2002.