

The Five Platonic Solids

Evan Finken

September 11, 2026

Abstract

This entry formalizes that there are exactly five Platonic solids: the tetrahedron, cube, octahedron, dodecahedron, and icosahedron. A Platonic solid is a convex polyhedron whose faces are congruent regular polygons, with the same number of edges meeting at each vertex. Euler’s polyhedron formula is used to show that there are at most five Platonic solids, and each of the five, defined as the convex hull of its vertices, is then shown to satisfy the required properties. The formal proof is computational (slow!) and is ported from John Harrison’s formalization in HOL Light [1]. Large language models were used to explain the intermediate steps of the HOL Light formalization and to write parts of the more complicated proofs.

Contents

1	Convexity	3
2	Determinants	7
3	Polytopes	8
3.1	Invariance under isometries	8
3.2	Distinctness of facets	10
3.3	Ridges lie in two facets	12
3.4	Lower bounds on face counts	22
3.4.1	Vertices	22
3.4.2	Facets	23
3.5	Double counting	26
3.5.1	$mF = 2E$	27
3.5.2	$nV = 2E$	28
3.6	Face containment	30
4	Computational helpers	31
4.1	Operations in $\mathbb{Q}(\sqrt{5})$	31
4.2	Operations in $\mathbb{R}^3$	35
4.3	Computation of faces	36

4.3.1	2-faces	36
4.3.2	1-faces	44
4.4	Full-dimensionality	53
4.5	Relations between faces of different dimensions	54
4.6	Congruence	56
4.6.1	Congruence of segments	57
4.6.2	Congruence of 2-faces	58
4.7	Equiangularity	71
5	Tetrahedron	74
5.1	Definition	74
5.2	Facets, edges, and vertices	75
5.3	Regularity	78
6	Cube	80
6.1	Definition	80
6.2	Facets, edges, and vertices	81
6.3	Regularity	86
7	Octahedron	88
7.1	Definition	88
7.2	Facets, edges, and vertices	89
7.3	Regularity	94
8	Dodecahedron	96
8.1	Definition	96
8.2	Facets, edges, and vertices	98
8.3	Regularity	108
9	Icosahedron	114
9.1	Definition	114
9.2	Facets, edges, and vertices	115
9.3	Regularity	128
10	Main result	132
10.1	Possible Schläfli symbols	132
10.2	Exactly five Platonic solids	135

1 Convexity

theory *Convex-Euclidean-Space-More*
imports *HOL-Analysis.Starlike*
begin

lemma *connected-Int-rel-frontier*:

assumes *connected* S
and $S \subseteq \text{affine hull } T$
and $S \cap T \neq \{\}$
and $S - T \neq \{\}$
shows $S \cap \text{rel-frontier } T \neq \{\}$

proof

assume *: $S \cap \text{rel-frontier } T = \{\}$
let ?E1 = $S \cap \text{rel-interior } T$
let ?E2 = $S - \text{closure } T$
have *openin* (*top-of-set* S) ?E1
by (*meson* *assms*(2) *openin-Int* *openin-rel-interior*
openin-subtopology-Int-subset *openin-subtopology-self*)
moreover **have** *openin* (*top-of-set* S) ?E2
by (*meson* *closed-closedin* *closed-closure* *closedin-self*
closedin-subtopology-refl *openin-subtopology-diff-closed*)
moreover **have** $S \subseteq ?E1 \cup ?E2$
using * *rel-frontier-def* **by** *fastforce*
moreover **have** $?E1 \cap ?E2 = \{\}$
using *rel-interior-subset-closure* **by** *fastforce*
moreover **have** $?E1 \neq \{\}$
using *assms*(3) *calculation*(3) *closure-subset* **by** *fastforce*
moreover **have** $?E2 \neq \{\}$
using *assms*(4) *calculation*(3) *rel-interior-subset* **by** *fastforce*
ultimately show *False*
using *connected-openin*[*of* S] *assms*(1) **by** *blast*

qed

end

theory *Convex-More*
imports *HOL-Analysis.Starlike*
begin

lemma *convex-hull-union-explicit*:

fixes $S T :: 'a::\text{euclidean-space set}$
assumes *convex* S
and *convex* T
shows *convex hull* ($S \cup T$) =
 $S \cup T \cup$
 $\{(1 - u) *_{\mathbb{R}} x + u *_{\mathbb{R}} y \mid x y u. x \in S \wedge y \in T \wedge 0 \leq u \wedge u \leq 1\}$
(is ?lhs = ?rhs)

proof

show ?lhs $\subseteq$?rhs

```

proof
  fix  $y$ 
  assume  $y \in \text{convex hull } (S \cup T)$ 
  then obtain  $U f$  where  $U\text{-facts}$ :
     $\text{finite } U \ U \subseteq (S \cup T) \ (\forall x \in U. \ 0 \leq f x) \ \text{sum } f \ U = 1 \ (\sum_{v \in U} f v *_{\mathbb{R}} v) = y$ 
    unfolding  $\text{convex-hull-explicit}$  by  $\text{auto}$ 
  have  $U\text{-partition}$ :  $U = (U \cap S) \cup (U - S)$ 
    by  $\text{auto}$ 
  have  $\text{disjoint}$ :  $(U \cap S) \cap (U - S) = \{\}$ 
    by  $\text{auto}$ 
  have  $\text{finite } (U \cap S) \ \text{finite } (U - S)$ 
    using  $U\text{-facts}(1)$  by  $\text{auto}$ 

  let  $?a = \text{sum } f \ (U \cap S)$ 
  let  $?b = \text{sum } f \ (U - S)$ 
  have  $?a + ?b = 1$ 
    using  $\text{sum.union-disjoint}[of \ U \cap S \ U - S \ f]$   $U\text{-facts}(1,4)$   $U\text{-partition disjoint}$ 
    by  $\text{auto}$ 
  have  $v\text{sum-eq}$ :  $y = (\sum_{v \in U \cap S} f v *_{\mathbb{R}} v) + (\sum_{v \in U - S} f v *_{\mathbb{R}} v)$ 
    using  $\text{sum.union-disjoint}[of \ U \cap S \ U - S \ \lambda v. f v *_{\mathbb{R}} v]$ 
    using  $U\text{-facts}(1,5)$   $U\text{-partition disjoint}$ 
    by  $\text{auto}$ 
  have  $0 \leq ?a \ 0 \leq ?b$ 
    using  $U\text{-facts}(3)$ 
    by  $(\text{auto intro: sum-nonneg})$ 

  consider
     $(a\text{-zero}) \ ?a = 0 \mid$ 
     $(b\text{-zero}) \ ?a \neq 0 \ ?b = 0 \mid$ 
     $(\text{both-pos}) \ ?a \neq 0 \ ?b \neq 0$ 
    by  $\text{auto}$ 
  then show  $y \in ?rhs$ 
  proof  $\text{cases}$ 
    case  $a\text{-zero}$ 
      then have  $\text{zero-on-}S$ :  $\forall x \in U \cap S. f x = 0$ 
        using  $U\text{-facts}(3)$   $\text{sum-nonneg-eq-0-iff}[of \ U \cap S \ f]$ 
        by  $(\text{simp add: } \langle \text{finite } (U \cap S) \rangle)$ 
      then have  $y = (\sum_{v \in U - S} f v *_{\mathbb{R}} v)$ 
        using  $v\text{sum-eq}$  by  $\text{simp}$ 
      moreover have  $\text{sum } f \ (U - S) = 1$ 
        using  $\langle ?a + ?b = 1 \rangle \ a\text{-zero}$ 
        by  $\text{simp}$ 
      moreover have  $U - S \subseteq T$ 
        using  $U\text{-facts}(2)$  by  $\text{auto}$ 
      moreover have  $\forall x \in U - S. 0 \leq f x$ 
        using  $U\text{-facts}(3)$  by  $\text{auto}$ 
      ultimately have  $y \in T$ 
        using  $\text{assms}(2) \ \langle \text{finite } (U - S) \rangle$  unfolding  $\text{convex-explicit}$ 
        by  $\text{blast}$ 

```

```

then show ?thesis by auto
next
case b-zero
then have zero-on-T:  $\forall x \in U - S. f\ x = 0$ 
  using U-facts(3) sum-nonneg-eq-0-iff[of U - S f]
  by (simp add:  $\langle \text{finite } (U - S) \rangle$ )
then have  $y = (\sum_{v \in U \cap S} f\ v *_{\mathbb{R}} v)$ 
  using vsum-eq by simp
moreover have  $\text{sum } f\ (U \cap S) = 1$ 
  using  $\langle ?a + ?b = 1 \rangle$  b-zero
  by simp
moreover have  $U \cap S \subseteq S$ 
  by auto
moreover have  $\forall x \in U \cap S. 0 \leq f\ x$ 
  using U-facts(3) by auto
ultimately have  $y \in S$ 
  using assms(1)  $\langle \text{finite } (U \cap S) \rangle$  unfolding convex-explicit
  by blast
then show ?thesis by auto
next
case both-pos
then have  $0 < ?a$   $0 < ?b$ 
  using  $\langle 0 \leq ?a \rangle \langle 0 \leq ?b \rangle$  by auto

let ?x =  $\sum_{v \in U \cap S} (f\ v / ?a) *_{\mathbb{R}} v$ 
let ?y' =  $\sum_{v \in U - S} (f\ v / ?b) *_{\mathbb{R}} v$ 

have  $\text{sum } (\lambda v. f\ v / ?a)\ (U \cap S) = 1$ 
  using  $\langle 0 < ?a \rangle$ 
  by (metis both-pos(1) div-self sum-divide-distrib)
moreover have  $\forall v \in U \cap S. 0 \leq f\ v / ?a$ 
  using U-facts(3)  $\langle 0 < ?a \rangle$  by auto
ultimately have  $?x \in S$ 
  using assms(1)  $\langle \text{finite } (U \cap S) \rangle$  unfolding convex-explicit
  by simp

have  $\text{sum } (\lambda v. f\ v / ?b)\ (U - S) = 1$ 
  using  $\langle 0 < ?b \rangle$ 
  by (metis both-pos(2) div-self sum-divide-distrib)
moreover have  $\forall v \in U - S. 0 \leq f\ v / ?b$ 
  using U-facts(3)  $\langle 0 < ?b \rangle$  by auto
moreover have  $U - S \subseteq T$ 
  using U-facts(2) by auto
ultimately have  $?y' \in T$ 
  using assms(2)  $\langle \text{finite } (U - S) \rangle$  unfolding convex-explicit
  by presburger

have  $(\sum_{v \in U \cap S} f\ v *_{\mathbb{R}} v) = ?a *_{\mathbb{R}} ?x$ 
  by (smt (verit) both-pos(1) eq-vector-fraction-iff scaleR-right.sum sum.cong)

```

```

moreover have  $(\sum_{v \in U - S} f \ v *_{\mathcal{R}} \ v) = ?b *_{\mathcal{R}} \ ?y'$ 
by (smt (verit) both-pos(2) eq-vector-fraction-iff scaleR-right.sum sum.cong)
ultimately have  $y = (1 - ?b) *_{\mathcal{R}} \ ?x + ?b *_{\mathcal{R}} \ ?y'$ 
using vsum-eq  $\langle ?a + ?b = 1 \rangle$  by auto

moreover have  $?b \leq 1$ 
using  $\langle 0 \leq ?a \rangle \langle 0 \leq ?b \rangle \langle ?a + ?b = 1 \rangle$ 
by auto

ultimately show ?thesis
using  $\langle 0 < ?b \rangle \langle ?x \in S \rangle \langle ?y' \in T \rangle$  by auto
qed
qed
next
have  $S \cup T \subseteq \text{convex hull } (S \cup T)$ 
using hull-subset by fast
then show ?rhs  $\subseteq$  ?lhs
by (smt (verit, ccfv-threshold) assms(1,2) convex-hull-union-two empty-iff
le-sup-iff
mem-Collect-eq subsetI)
qed

lemma convex-hull-union-nonempty-explicit:
fixes  $S \ T :: 'a::\text{euclidean-space set}$ 
assumes convex  $S$ 
and  $S \neq \{\}$ 
and convex  $T$ 
and  $T \neq \{\}$ 
shows  $\text{convex hull } (S \cup T) =$ 
 $\{(1 - u) *_{\mathcal{R}} \ x + u *_{\mathcal{R}} \ y \mid x \ y \ u. \ x \in S \wedge y \in T \wedge 0 \leq u \wedge u \leq 1\}$ 
(is - = ?rhs)
proof -
have  $S \subseteq ?rhs$ 
using assms(4) by fastforce
moreover have  $T \subseteq ?rhs$ 
by (smt (verit, ccfv-threshold) add-0 assms(2) equals0I mem-Collect-eq scaleR-one
scaleR-zero-left subset-iff)
ultimately show ?thesis
using convex-hull-union-explicit[OF assms(1,3)]
by blast
qed

lemma convex-hull-isometry:
assumes orthogonal-transformation  $f$ 
shows  $(\lambda x. c + f \ x) \ ` (\text{convex hull } S) = \text{convex hull } ((\lambda x. c + f \ x) \ ` S)$ 
using orthogonal-transformation-linear[OF assms] convex-hull-linear-image[of  $f$ 
 $S$ ]
convex-hull-translation[of  $c \ f \ ` S$ ]
by (metis image-image)

```

end

2 Determinants

theory *Determinants-More*
 imports *HOL-Analysis.Starlike*
 begin

lemma *hyperplane-subset-imp-scaled*:

fixes $a :: 'a::euclidean-space$

assumes $\{x. a \cdot x = b\} \subseteq \{x. a' \cdot x = b'\}$

and $a \neq 0$

and $a' \neq 0$

shows $\exists c. a' = c *_{\mathbb{R}} a \wedge b' = c * b \wedge c \neq 0$

proof –

let $?c = (a \cdot a') / (a \cdot a)$

let $?v = a' - ?c *_{\mathbb{R}} a$

have $a \cdot ?v = 0$

by (*simp add: inner-diff*)

let $?w = (b / (a \cdot a)) *_{\mathbb{R}} a$

have $a \cdot ?w = b$

using *assms(2)* by *simp*

then have $a' \cdot ?w = b'$

using *assms(1)* by *blast*

have $a \cdot (?w + ?v) = b$

using $\langle a \cdot ?w = b \rangle \langle a \cdot ?v = 0 \rangle$

by (*simp add: inner-add-right*)

then have $a' \cdot (?w + ?v) = b'$

using *assms(1)* by *blast*

then have $a' \cdot ?v = 0$

using $\langle a' \cdot ?w = b' \rangle$

by (*simp add: inner-add-right*)

have $?v \cdot ?v = 0$

using $\langle a' \cdot ?v = 0 \rangle \langle a \cdot ?v = 0 \rangle$

by (*simp add: inner-diff-left*)

then have $?v = 0$

by *simp*

then have $a' = ?c *_{\mathbb{R}} a$

by *auto*

moreover have $b' = ?c * b$

by (*metis (mono-tags, lifting) assms(1) Collect-mono-iff* $\langle a' = ?c *_{\mathbb{R}} a \rangle$ *hyperplane-eq-Ex*)

inner-scaleR-left assms(2))

moreover have $?c \neq 0$

using $\langle a' = ?c *_{\mathbb{R}} a \rangle$ *assms(3)* by *fastforce*

ultimately show *?thesis*

by *blast*

qed

```

lemma subset-hyperplanes:
  fixes  $a :: 'a::euclidean-space$ 
  shows  $(\{x. a \cdot x = b\} \subseteq \{x. a' \cdot x = b'\}) =$ 
     $(\{x. a \cdot x = b\} = \{\} \vee \{x. a' \cdot x = b'\} = UNIV \vee$ 
     $\{x. a \cdot x = b\} = \{x. a' \cdot x = b'\})$ 
  (is ?lhs = ?rhs)
proof
  assume *: ?lhs
  consider
     $(a\text{-zero})\ a = 0 \mid$ 
     $(a'\text{-zero})\ a' = 0 \mid$ 
     $(nonzero)\ a \neq 0\ a' \neq 0$ 
  by blast
  then show ?rhs
  proof (cases)
    case a-zero
    then show ?thesis using * by auto
  next
    case a'-zero
    then show ?thesis using * by force
  next
    case nonzero
    then obtain  $c$  where  $\langle a' = c *_R a \rangle \langle b' = c * b \rangle$ 
    using * hyperplane-subset-imp-scaled by blast
    then show ?thesis
    by (smt (verit, ccv-threshold) Collect-cong inner-scaleR-left
      nonzero(2) scaleR-eq-0-iff vector-space-over-itself.scale-left-imp-eq)
  qed
next
  assume ?rhs
  then show ?lhs by blast
qed

end

```

3 Polytopes

```

theory Polytope-More
  imports
    Convex-Euclidean-Space-More
    HOL-Analysis.Analysis
begin

```

3.1 Invariance under isometries

```

lemma facet-of-translation-eq:
   $((+) a \text{ ' } T \text{ facet-of } (+) a \text{ ' } S) = (T \text{ facet-of } S)$ 
  by (simp add: aff-dim-translation-eq facet-of-def)

```

lemma *facets-of-translation*:
shows $\{F. F \text{ facet-of } (+) a \text{ ' } S\} = (') ((+) a) \text{ ' } \{F. F \text{ facet-of } S\}$
proof –
have $\bigwedge F. F \text{ facet-of } (+) a \text{ ' } S \implies \exists G. G \text{ facet-of } S \wedge F = (+) a \text{ ' } G$
by (*metis facet-of-imp-subset facet-of-translation-eq subset-imageE*)
then show *?thesis*
by (*auto simp: image-iff facet-of-translation-eq*)
qed

lemma *polyhedron-translation-eq*:
shows $\text{polyhedron } S = \text{polyhedron } ((+) x \text{ ' } S)$
proof –
have *translation*: $\bigwedge S x. \text{polyhedron } S \implies \text{polyhedron } ((+) x \text{ ' } S)$
proof –
fix $S::'b \text{ set}$ **and** $x::'b$
assume $\text{polyhedron } S$
then obtain F **where** $F\text{-props}$:
 $\text{finite } F$
 $S = \bigcap F$
 $\bigwedge h. h \in F \implies \exists a b. a \neq 0 \wedge h = \{y. a \cdot y \leq b\}$
unfolding *polyhedron-def* **by** *force*
let $?F' = (') ((+) x) \text{ ' } F$
have *finite* $?F'$
using $F\text{-props}(1)$ **by** *blast*
moreover have $((+) x \text{ ' } S) = \bigcap ?F'$
by (*metis (no-types, lifting) F-props(1,2) Inf-fin.hom-commute*
 $\text{Inf-fin-Inf Inf-top-conv}(2) \text{ calculation empty-iff empty-is-image}$
 $\text{surj-plus translation-Int}$)
moreover have $\bigwedge h'. h' \in ?F' \implies \exists a b. a \neq 0 \wedge h' = \{y. a \cdot y \leq b\}$
proof –
fix h'
assume $h' \in ?F'$
then obtain $h a b$ **where**
 $h'\text{-eq}: h' = ((+) x) \text{ ' } h$ **and**
 $h\text{-props}: h \in F a \neq 0 h = \{y. a \cdot y \leq b\}$
using $F\text{-props}(3)$ **by** *blast*
have $h' = \{y + x \mid y. a \cdot y \leq b\}$
using $h'\text{-eq } h\text{-props}(3)$ **by** *auto*
also have $\dots = \{y. a \cdot (y - x) \leq b\}$
by *force*
finally have $h' = \{y. a \cdot y \leq b + a \cdot x\}$
by (*auto simp: inner-diff-right*)
then show $\exists a b. a \neq 0 \wedge h' = \{y. a \cdot y \leq b\}$
using $h\text{-props}(2)$ **by** *blast*
qed

ultimately show $\text{polyhedron } ((+) x \text{ ' } S)$
unfolding *polyhedron-def* **by** *meson*

qed
 show ?thesis
 using translation[of S x] translation[of $((+) x \text{ ' } S) - x$]
 by (metis translation-galois)
 qed

lemma face-of-isometry-eq:
 assumes orthogonal-transformation f
 shows $((\lambda x. c + f x) \text{ ' } T \text{ face-of } (\lambda x. c + f x) \text{ ' } S) \longleftrightarrow T \text{ face-of } S$
 by (metis assms face-of-translation-eq face-of-linear-image image-image
 orthogonal-transformation-inj orthogonal-transformation-linear)

lemma extreme-point-of-isometry-eq:
 assumes orthogonal-transformation f
 shows $(c + f T \text{ extreme-point-of } (\lambda x. c + f x) \text{ ' } S) \longleftrightarrow T \text{ extreme-point-of } S$
 using face-of-isometry-eq[OF assms, of $c \{T\} S$]
 by (simp add: face-of-singleton)

lemma aff-dim-isometry-eq:
 assumes orthogonal-transformation f
 shows $\text{aff-dim } S = \text{aff-dim } ((\lambda x. c + f x) \text{ ' } S)$
 by (metis assms aff-dim-injective-linear-image aff-dim-translation-eq image-image
 orthogonal-transformation-def orthogonal-transformation-inj)

3.2 Distinctness of facets

lemma facets-of-polyhedron-explicit-distinct:
 fixes $S :: 'a :: \text{euclidean-space set}$
 assumes finite: finite F
 and seq: $S = \text{affine hull } S \cap \bigcap F$
 and faceq: $\bigwedge h. h \in F \implies a h \neq 0 \wedge h = \{x. a h \cdot x \leq b h\}$
 and psub: $\bigwedge F'. F' \subset F \implies S \subset \text{affine hull } S \cap \bigcap F'$
 and h1-in: $h1 \in F$
 and h2-in: $h2 \in F$
 and eq: $S \cap \{x. a h1 \cdot x = b h1\} = S \cap \{x. a h2 \cdot x = b h2\}$
 shows $h1 = h2$
 proof (rule ccontr)
 assume *: $h1 \neq h2$
 have $S \neq \{\}$
 using h1-in psub by force
 have polyhedron S
 by (meson polyhedron-Int-affine[of S] seq faceq finite)
 have rel-interior $S \neq \{\}$
 using rel-interior-polyhedron-explicit
 by (simp add: $\langle S \neq \{\} \rangle \langle \text{polyhedron } S \rangle \text{polyhedron-imp-convex}$
 rel-interior-eq-empty)
 then obtain z where z -props: $z \in S \wedge h. h \in F \implies a h \cdot z < b h$
 using rel-interior-polyhedron-explicit[OF finite seq] faceq psub
 by blast

```

have  $S \subset \text{affine hull } S \cap \bigcap (F - \{h2\})$ 
  using  $h2\text{-in psub}[of F - \{h2\}]$ 
  by blast
then obtain  $x$  where  $x\text{-props: } x \in \text{affine hull } S \wedge x \in \bigcap (F - \{h2\}) \wedge x \notin S$ 
  by auto
then have  $a \cdot h2 \cdot x > b \cdot h2$ 
  using Diff-iff empty-iff faceq insert-iff seq
  by fastforce
have  $\bigwedge h. h \in (F - \{h2\}) \implies a \cdot h \cdot x \leq b \cdot h$ 
  using  $x\text{-props}(2)$  faceq by auto
have  $\text{closed-segment } x \ z \cap \text{rel-frontier } S \neq \{\}$ 
  by (metis Diff-iff Int-iff closed-segment-subset connected-Int-rel-frontier
    connected-segment convex-affine-hull empty-iff ends-in-segment(1,2) seq
     $x\text{-props}(1,3)$   $z\text{-props}(1)$ )
then have  $\text{closed-segment } x \ z \cap (S - \text{rel-interior } S) \neq \{\}$ 
  by (simp add: polyhedron S rel-boundary-of-polyhedron
    rel-frontier-of-polyhedron)
then obtain  $y$  where  $y\text{-props:}$ 
   $y \in \text{closed-segment } x \ z$ 
   $y \in S - \text{rel-interior } S$ 
   $y \in S$ 
  by blast
have  $S - \text{rel-interior } S = \bigcup \{S \cap \{x. a \cdot h \cdot x = b \cdot h\} \mid h. h \in F\}$ 
  using facet-of-polyhedron-explicit[OF finite seq faceq psub]
    rel-boundary-of-polyhedron[OF polyhedron S]
  by auto
then obtain  $h$  where  $h\text{-props: } h \in F \wedge a \cdot h \cdot y = b \cdot h$ 
  using  $y\text{-props}(2)$  by auto
have  $\exists k. k \in F \wedge k \neq h2 \wedge a \cdot k \cdot y = b \cdot k$ 
proof (cases h = h2)
  case True
    then have  $y \in S \cap \{x. a \cdot h2 \cdot x = b \cdot h2\}$ 
    using  $h\text{-props}(2)$   $y\text{-props}(3)$  by blast
    then have  $a \cdot h1 \cdot y = b \cdot h1$ 
    using eq by blast
    then show ?thesis
    using h1-in * by blast
  next
    case False
    then show ?thesis
    using  $h\text{-props}$  by blast
qed
then obtain  $k$  where  $k\text{-props: } k \in F \wedge k \neq h2 \wedge a \cdot k \cdot y = b \cdot k$ 
  by blast
have  $y \in \text{open-segment } x \ z$ 
  by (metis DiffI polyhedron S rel-boundary-of-polyhedron
    all-not-in-conv insertE open-segment-def order-less-irrefl
     $x\text{-props}(3)$   $y\text{-props}(1,3)$   $z\text{-props}(2)$ )
then obtain  $u$  where  $u\text{-props: } y = (1 - u) \cdot x + u \cdot z \wedge 0 < u < 1$ 

```

```

    using in-segment(2)[of y x z] by blast
  have  $(1 - u) * (a \cdot k \cdot x) + u * (a \cdot k \cdot z) = b \cdot k$ 
    using k-props(3) unfolding u-props(1)
    by (simp add: inner-right-distrib)
  moreover have  $(1 - u) * (a \cdot k \cdot x) \leq (1 - u) * b \cdot k$ 
    by (simp add: ‹ $\bigwedge ha. ha \in F - \{h2\} \implies a \cdot ha \cdot x \leq b \cdot ha$ › k-props(1,2)
    u-props(3))
  moreover have  $u * (a \cdot k \cdot z) < u * b \cdot k$ 
    by (simp add: k-props(1) u-props(2) z-props(2))
  ultimately show False
    by argo
qed

```

3.3 Ridges lie in two facets

```

lemma polyhedron-ridge-two-facets-0:
  fixes p :: 'a :: euclidean-space set
  assumes polyhedron p
    and r face-of p
    and 0 ∈ r
    and aff-dim r = aff-dim p - 2
  shows  $\exists f1 f2. f1 \text{ face-of } p \wedge \text{aff-dim } f1 = \text{aff-dim } p - 1 \wedge$ 
     $f2 \text{ face-of } p \wedge \text{aff-dim } f2 = \text{aff-dim } p - 1 \wedge$ 
     $f1 \neq f2 \wedge r \subseteq f1 \wedge r \subseteq f2 \wedge f1 \cap f2 = r \wedge$ 
     $(\forall f. f \text{ face-of } p \wedge \text{aff-dim } f = \text{aff-dim } p - 1 \wedge r \subseteq f$ 
       $\longrightarrow f = f1 \vee f = f2)$ 

proof -
  let ?S = {f. f face-of p ∧ aff-dim f = aff-dim p - 1 ∧ r ⊆ f}
  have finite ?S
    by (simp add: asms(1) finite-polyhedron-faces)
  have r = ⋂ {f. f facet-of p ∧ r ⊆ f}
    using face-of-polyhedron[of p r] asms by fastforce
  then have r-eq: r = ⋂ ?S
    using asms(3) facet-of-def by blast
  have 0 ∈ p
    using asms(2,3) face-of-imp-subset by auto

  have card ?S = 0 ∨ card ?S = 1 ∨ card ?S = 2 ∨ card ?S ≥ 3
    by linarith
  then consider
    (card0) card ?S = 0 |
    (card1) card ?S = 1 |
    (card2) card ?S = 2 |
    (card3) card ?S ≥ 3
    by argo

  then show ?thesis
  proof (cases)
    case card0

```

then show *?thesis*
by (*smt* (*verit*, *best*) *Inf-empty* $\langle \text{finite } ?S \rangle$ *aff-dim-UNIV*
aff-dim-le-DIM *assms*(4) *card-gt-0-iff* *nless-le* *r-eq*)

next

case *card1*
then obtain *f1* **where**
S-eq: $?S = \{f1\}$ **and**
aff-dim-f1: $\text{aff-dim } f1 = \text{aff-dim } p - 1$
by (*metis* (*mono-tags*, *lifting*) *all-not-in-conv* *card-1-singletonE*
insert-not-empty *mem-Collect-eq* *singletonD*)
show *?thesis*
using *r-eq* *aff-dim-f1* *assms*(4) **unfolding** *S-eq*
by *simp*

next

case *card2*
then obtain *f1 f2* **where**
S-eq: $?S = \{f1, f2\}$ **and**
f1 face-of p $\text{aff-dim } f1 = \text{aff-dim } p - 1$ $r \subseteq f1$ **and**
f2 face-of p $\text{aff-dim } f2 = \text{aff-dim } p - 1$ $r \subseteq f2$ **and**
 $f1 \neq f2$
by (*smt* (*verit*) *card-2-iff* *insert-iff* *mem-Collect-eq*)
moreover have $f1 \cap f2 = r$
using *S-eq* *r-eq* **by** *auto*
moreover have $(\forall f. f \text{ face-of } p \wedge \text{aff-dim } f = \text{aff-dim } p - 1 \wedge r \subseteq f$
 $\longrightarrow f = f1 \vee f = f2)$
using *S-eq* **by** *auto*
ultimately show *?thesis*
by *blast*

next

case *card3*
then have $\text{card } ?S \geq \text{Suc } (\text{Suc } (\text{Suc } 0))$
by *presburger*
then obtain *f1 f2 f3 S2* **where**
S-eq: $?S = \text{insert } f3 (\text{insert } f2 (\text{insert } f1 S2))$ **and**
fs-neg: $f1 \neq f2$ $f1 \neq f3$ $f2 \neq f3$
by (*smt* (*verit*) *card-le-Suc-iff* *insertCI*)
have *f1-props*: $f1 \text{ face-of } p$ $\text{aff-dim } f1 = \text{aff-dim } p - 1$ $r \subseteq f1$ **and**
f2-props: $f2 \text{ face-of } p$ $\text{aff-dim } f2 = \text{aff-dim } p - 1$ $r \subseteq f2$ **and**
f3-props: $f3 \text{ face-of } p$ $\text{aff-dim } f3 = \text{aff-dim } p - 1$ $r \subseteq f3$
using *S-eq*
by *auto*

obtain *F a b* **where**

finite F **and**
p-eq: $p = \text{affine hull } p \cap \bigcap F$ **and**
halfspaces: $\bigwedge h. h \in F \implies a \cdot h \neq 0 \wedge h = \{x. a \cdot h \cdot x \leq b\} \wedge$

$(\forall x \in \text{affine hull } p. x + a \cdot h \in \text{affine hull } p)$ **and**
minimal: $\bigwedge F'. F' \subset F \longrightarrow p \subset \text{affine hull } p \cap \bigcap F'$
using *polyhedron-Int-affine-parallel-minimal*[of p] *assms*(1)
by *metis*
obtain $h1$ **where** $h1 \in F$ $a \cdot h1 \neq 0$ **and** $f1\text{-eq}$: $f1 = p \cap \{x. a \cdot h1 \cdot x = b \cdot h1\}$
using *facet-of-polyhedron-explicit*[OF $\langle \text{finite } F \rangle$ $p\text{-eq}$, of a b $f1$]
halfspaces minimal f1-props unfolding facet-of-def
using *assms*(3) **by** *blast*
obtain $h2$ **where** $h2 \in F$ $a \cdot h2 \neq 0$ **and** $f2\text{-eq}$: $f2 = p \cap \{x. a \cdot h2 \cdot x = b \cdot h2\}$
using *facet-of-polyhedron-explicit*[OF $\langle \text{finite } F \rangle$ $p\text{-eq}$, of a b $f2$]
halfspaces minimal f2-props unfolding facet-of-def
using *assms*(3) **by** *blast*
obtain $h3$ **where** $h3 \in F$ $a \cdot h3 \neq 0$ **and** $f3\text{-eq}$: $f3 = p \cap \{x. a \cdot h3 \cdot x = b \cdot h3\}$
using *facet-of-polyhedron-explicit*[OF $\langle \text{finite } F \rangle$ $p\text{-eq}$, of a b $f3$]
halfspaces minimal f3-props unfolding facet-of-def
using *assms*(3) **by** *blast*

have *all-x-in-r*: $\bigwedge x. x \in r \implies a \cdot h1 \cdot x = b \cdot h1 \wedge a \cdot h2 \cdot x = b \cdot h2 \wedge a \cdot h3 \cdot x = b \cdot h3$
using *f1-props*(3) *f2-props*(3) *f3-props*(3) *f1-eq* *f2-eq* *f3-eq*
by *blast*
then have *b-hs-eq-0*: $b \cdot h1 = 0 \wedge b \cdot h2 = 0 \wedge b \cdot h3 = 0$
using $\langle 0 \in r \rangle$ **by** *force+*
then have *all-x-in-r-0*: $\bigwedge x. x \in r \implies a \cdot h1 \cdot x = 0 \wedge a \cdot h2 \cdot x = 0 \wedge a \cdot h3 \cdot x = 0$
using *all-x-in-r* **by** *auto*
have *a-hs-neq*: $a \cdot h1 \neq 0 \wedge a \cdot h2 \neq 0 \wedge a \cdot h3 \neq 0$
using *b-hs-eq-0* *f1-eq* *f2-eq* *f3-eq* *fs-neq* **by** *force+*
have *hs-neq*: $h1 \neq h2 \wedge h1 \neq h3 \wedge h2 \neq h3$
using *fs-neq* *f1-eq* *f2-eq* *f3-eq* **by** *auto*

have *aff-dim* $r = \text{dim } r$
using *aff-dim-zero*[OF *hull-inc*[OF $\langle 0 \in r \rangle$]] .
moreover have *aff-dim* $p = \text{dim } p$
using *aff-dim-zero*[OF *hull-inc*[OF $\langle 0 \in p \rangle$]] .
ultimately have $\text{dim } r = \text{dim } p - 2$
using *assms*(4) **by** *presburger*

have *not-contained-xy-z*:
 $\neg (\{x. a \cdot hx \cdot x \leq 0\} \cap \{x. a \cdot hy \cdot x \leq 0\} \subseteq \{x. a \cdot hz \cdot x \leq 0\})$
if $hx \in F$ $hy \in F$ $hz \in F$ $b \cdot hx = 0 \wedge b \cdot hy = 0 \wedge b \cdot hz = 0$ $hx \neq hz$ $hy \neq hz$ **for** hx
 hy hz
proof
assume $\{x. a \cdot hx \cdot x \leq 0\} \cap \{x. a \cdot hy \cdot x \leq 0\} \subseteq \{x. a \cdot hz \cdot x \leq 0\}$
then have *: $hx \cap hy \subseteq hz$
using *that halfspaces*
by *metis*
have $F - \{hz\} \subset F$
using $\langle hz \in F \rangle$ $\langle \text{finite } F \rangle$ **by** *blast*

```

then have  $p$ -psubset:  $p \subset \text{affine hull } p \cap \bigcap (F - \{hz\})$ 
  using minimal by presburger
have  $hx \in F - \{hz\}$   $hy \in F - \{hz\}$ 
  using that by auto
then show False
  using *  $p$ -eq  $p$ -psubset by blast
qed
have not-contained-12-3:
   $\neg (\{x. a \ h1 \cdot x \leq 0\} \cap \{x. a \ h2 \cdot x \leq 0\} \subseteq \{x. a \ h3 \cdot x \leq 0\})$ 
  using not-contained-xy-z[ $OF \ \langle h1 \in F \rangle \ \langle h2 \in F \rangle \ \langle h3 \in F \rangle$ ] b-hs-eq-0 hs-neq
  by blast
have not-contained-13-2:
   $\neg (\{x. a \ h1 \cdot x \leq 0\} \cap \{x. a \ h3 \cdot x \leq 0\} \subseteq \{x. a \ h2 \cdot x \leq 0\})$ 
  using not-contained-xy-z[ $OF \ \langle h1 \in F \rangle \ \langle h3 \in F \rangle \ \langle h2 \in F \rangle$ ] b-hs-eq-0 hs-neq
  by blast
have not-contained-23-1:
   $\neg (\{x. a \ h2 \cdot x \leq 0\} \cap \{x. a \ h3 \cdot x \leq 0\} \subseteq \{x. a \ h1 \cdot x \leq 0\})$ 
  using not-contained-xy-z[ $OF \ \langle h2 \in F \rangle \ \langle h3 \in F \rangle \ \langle h1 \in F \rangle$ ] b-hs-eq-0 hs-neq
  by blast

obtain  $w$  where  $w \in \text{rel-interior } p$ 
  using  $\langle 0 \in p \rangle$  polyhedron-imp-convex rel-interior-eq-empty assms(1)
  by auto
moreover have  $\text{rel-interior } p = \{x \in p. \forall h \in F. a \ h \cdot x < b \ h\}$ 
  using rel-interior-polyhedron-explicit[ $OF \ \langle \text{finite } F \rangle \ p$ -eq, of a b]
  halfspaces minimal
  by blast
ultimately have  $w$ -props:  $a \ h1 \cdot w < 0 \ a \ h2 \cdot w < 0 \ a \ h3 \cdot w < 0$ 
  using b-hs-eq-0  $\langle h1 \in F \rangle \ \langle h2 \in F \rangle \ \langle h3 \in F \rangle$ 
  by auto

have  $r \subseteq \text{span } p$ 
  using face-of-imp-subset[ $OF \ \text{assms}(2)$ ] span-superset[of p]
  by blast
have  $a \ h1 \in \text{span } p \ a \ h2 \in \text{span } p \ a \ h3 \in \text{span } p$ 
  using halfspaces  $\langle h1 \in F \rangle \ \langle h2 \in F \rangle \ \langle h3 \in F \rangle \ \langle 0 \in p \rangle$ 
  by (metis add-0 affine-hull-span-0 hull-inc)+
then have  $\text{span } \{a \ h1, a \ h2, a \ h3\} \subseteq \text{span } p$ 
  using subspace-span[of p]
  by (simp add: span-minimal)
then have  $r \cup \text{span } \{a \ h1, a \ h2, a \ h3\} \subseteq \text{span } p$ 
  using  $\langle r \subseteq \text{span } p \rangle$  subspace-span[of p]
  by (simp add: subspace-sum-minimal)
have  $\dim (\text{span } p) = \dim p$ 
  by simp
have  $(\bigwedge x \ y. x \in r \implies y \in \text{span } \{a \ h1, a \ h2, a \ h3\} \implies x \cdot y = 0)$ 
proof -
  fix  $x \ y$ 
  assume  $x$ -in- $r$ :  $x \in r$ 

```

```

assume  $y \in \text{span } \{a \ h1, a \ h2, a \ h3\}$ 
then show  $x \cdot y = 0$ 
proof (induct y rule: span-induct)
  case base
  then show ?case
    using subspace-hyperplane by blast
  next
  case (step y)
  then show ?case
    using all-x-in-r-0[OF x-in-r]
    by (metis inner-commute empty-iff insertE)
  qed
qed
then have  $\dim (r \cup \text{span } \{a \ h1, a \ h2, a \ h3\}) = \dim r + \dim (\text{span } \{a \ h1, a \ h2, a \ h3\})$ 
  using dim-orthogonal-sum[of r span {a h1, a h2, a h3}]
  by blast
moreover have  $\dim (r \cup \text{span } \{a \ h1, a \ h2, a \ h3\}) \leq \dim p$ 
  using dim-subset[OF  $\langle r \cup \text{span } \{a \ h1, a \ h2, a \ h3\} \subseteq \text{span } p \rangle$ ]
  by auto
ultimately have  $\dim (\text{span } \{a \ h1, a \ h2, a \ h3\}) \leq 2$ 
  using  $\langle \dim r = \dim p - 2 \rangle$ 
  by linarith

have independent {a h1, a h2, a h3}
proof
  assume dependent {a h1, a h2, a h3}
  have finite {a h1, a h2, a h3} by simp
  obtain c where c-sum:  $(\sum_{v \in \{a \ h1, a \ h2, a \ h3\}} c \ v \ *_R \ v) = 0$ 
    and nonzero:  $\exists v \in \{a \ h1, a \ h2, a \ h3\}. c \ v \neq 0$ 
    using  $\langle \text{dependent } \{a \ h1, a \ h2, a \ h3\} \rangle$  dependent-finite[OF  $\langle \text{finite } \{a \ h1, a \ h2, a \ h3\} \rangle$ ]
    by auto

  let ?c1 =  $c \ (a \ h1)$  and ?c2 =  $c \ (a \ h2)$  and ?c3 =  $c \ (a \ h3)$ 
  have sum-0:  $?c1 \ *_R \ a \ h1 + ?c2 \ *_R \ a \ h2 + ?c3 \ *_R \ a \ h3 = 0$ 
    using c-sum a-hs-neq
    by (simp add: group-cancel.add1)
  have ?c1  $\neq 0$ 
proof
  assume ?c1 = 0
  then have  $?c2 \ *_R \ a \ h2 + ?c3 \ *_R \ a \ h3 = 0$  using sum-0 by simp
  then have ?c2  $\neq 0$  ?c3  $\neq 0$ 
    using nonzero  $\langle ?c1 = 0 \rangle$ 
    using  $\langle a \ h2 \neq 0 \rangle \langle a \ h3 \neq 0 \rangle$  by auto

  define k where  $k = - ( ?c3 \ / \ ?c2 )$ 
  have  $a \ h2 = k \ *_R \ a \ h3$ 
    using  $\langle ?c2 \ *_R \ a \ h2 + ?c3 \ *_R \ a \ h3 = 0 \rangle \langle ?c2 \neq 0 \rangle$ 

```

```

    unfolding k-def
    by (smt (verit, ccfv-SIG) eq-vector-fraction-iff neg-eq-iff-add-eq-0 scaleR-left.minus
scaleR-minus-right)

show False
proof (cases k > 0)
  case True
  then have {x. a h3 · x ≤ 0} = {x. a h2 · x ≤ 0}
    unfolding ⟨a h2 = k *R a h3⟩ by (auto simp: mult-le-0-iff)
  then have {x. a h1 · x ≤ 0} ∩ {x. a h2 · x ≤ 0} ⊆ {x. a h3 · x ≤ 0}
    by blast
  then show False using not-contained-12-3 by simp
next
  case False
  then have k < 0
    using ⟨a h2 = k *R a h3⟩ a-hs-neg(3) ⟨a h2 ≠ 0⟩ by fastforce
  have a h2 · w = k * (a h3 · w) using ⟨a h2 = k *R a h3⟩ ⟨a h2 ≠ 0⟩
by auto
    moreover have a h2 · w < 0 using w-props(2) by simp
    moreover have k * (a h3 · w) > 0 using ⟨k < 0⟩ w-props(3) mult-neg-neg
by blast
    ultimately show False by simp
qed
qed

define u where u = -(?c2 / ?c1)
define v where v = -(?c3 / ?c1)
have ?c1 *R a h1 = -?c2 *R a h2 - ?c3 *R a h3
  by (metis (no-types, lifting) sum-0 add.left-commute add.right-neutral
add-minus-cancel scaleR-left.minus
uminus-add-conv-diff)
then have (1 / ?c1) *R (?c1 *R a h1) = (1 / ?c1) *R (-?c2 *R a h2 - ?c3
*R a h3)
  by auto
then have a1-eq: a h1 = u *R a h2 + v *R a h3
  using ⟨?c1 ≠ 0⟩ unfolding u-def v-def
  by (simp add: scaleR-right-diff-distrib)

have u ≠ 0
proof
  assume u = 0
  then have a h1 = v *R a h3
    using a1-eq by simp
  show False
proof (cases v > 0)
  case True
  then have {x. a h3 · x ≤ 0} = {x. a h1 · x ≤ 0}
    unfolding ⟨a h1 = v *R a h3⟩
    by (auto simp: mult-le-0-iff)

```

```

then have  $\{x. a \cdot h2 \cdot x \leq 0\} \cap \{x. a \cdot h3 \cdot x \leq 0\} \subseteq \{x. a \cdot h1 \cdot x \leq 0\}$ 
  by blast
then show False
  using not-contained-23-1
  by simp
next
case False
then have  $v < 0$ 
  using  $\langle a \cdot h1 = v *_R a \cdot h3 \rangle$  a-hs-neq(2)  $\langle a \cdot h1 \neq 0 \rangle$ 
  by auto
have  $a \cdot h1 \cdot w = v * (a \cdot h3 \cdot w)$ 
  using  $\langle a \cdot h1 = v *_R a \cdot h3 \rangle$ 
  by simp
moreover have  $a \cdot h1 \cdot w < 0$ 
  using w-props(1) by simp
moreover have  $v * (a \cdot h3 \cdot w) > 0$ 
  using  $\langle v < 0 \rangle$  w-props(3) mult-neg-neg
  by blast
ultimately show False by simp
qed
qed

have  $v \neq 0$ 
proof
assume  $v = 0$ 
then have  $a \cdot h1 = u *_R a \cdot h2$ 
  using a1-eq by simp
show False
proof (cases  $u > 0$ )
case True
then have  $\{x. a \cdot h2 \cdot x \leq 0\} = \{x. a \cdot h1 \cdot x \leq 0\}$ 
  unfolding  $\langle a \cdot h1 = u *_R a \cdot h2 \rangle$ 
  by (auto simp: mult-le-0-iff)
then have  $\{x. a \cdot h2 \cdot x \leq 0\} \cap \{x. a \cdot h3 \cdot x \leq 0\} \subseteq \{x. a \cdot h1 \cdot x \leq 0\}$ 
  by blast
then show False
  using not-contained-23-1
  by simp
next
case False
then have  $u < 0$ 
  using  $\langle a \cdot h1 = u *_R a \cdot h2 \rangle$  a-hs-neq(1)  $\langle a \cdot h1 \neq 0 \rangle$ 
  by auto
have  $a \cdot h1 \cdot w = u * (a \cdot h2 \cdot w)$ 
  using  $\langle a \cdot h1 = u *_R a \cdot h2 \rangle$ 
  by simp
moreover have  $a \cdot h1 \cdot w < 0$ 
  using w-props(1) by simp
moreover have  $u * (a \cdot h2 \cdot w) > 0$ 

```

```

    using ⟨u < 0⟩ w-props(2) mult-neg-neg
    by blast
  ultimately show False by simp
qed
qed

consider (pos-pos) u > 0 v > 0 |
  (pos-neg) u > 0 v < 0 |
  (neg-pos) u < 0 v > 0 |
  (neg-neg) u < 0 v < 0
  using ⟨u ≠ 0⟩ ⟨v ≠ 0⟩ by linarith
then show False
proof cases
  case pos-pos
  have {x. a h2 · x ≤ 0} ∩ {x. a h3 · x ≤ 0} ⊆ {x. a h1 · x ≤ 0}
  proof
    fix x
    assume x-in: x ∈ {x. a h2 · x ≤ 0} ∩ {x. a h3 · x ≤ 0}
    have a h1 · x = u * (a h2 · x) + v * (a h3 · x)
      using a1-eq by (simp add: inner-add)
    also have ... ≤ 0
      using x-in pos-pos
      by (simp add: add-decreasing2 split-mult-neg-le)
    finally show x ∈ {x. a h1 · x ≤ 0}
      by blast
  qed
then show False
  using not-contained-23-1 by simp
next
  case pos-neg
  have u *R a h2 = a h1 + (-v) *R a h3
    using a1-eq
    by (simp add: algebra-simps)
  have {x. a h1 · x ≤ 0} ∩ {x. a h3 · x ≤ 0} ⊆ {x. a h2 · x ≤ 0}
  proof
    fix x assume x-in: x ∈ {x. a h1 · x ≤ 0} ∩ {x. a h3 · x ≤ 0}
    have u * (a h2 · x) = (a h1 · x) + (-v) * (a h3 · x)
      using ⟨u *R a h2 = a h1 + (-v) *R a h3⟩
      by (metis inner-left-distrib inner-scaleR-left)
    also have ... ≤ 0
      using x-in pos-neg
      by (smt (verit, best) IntE mem-Collect-eq mult-le-0-iff)
    finally have u * (a h2 · x) ≤ 0 .
    then show x ∈ {x. a h2 · x ≤ 0}
      using ⟨u > 0⟩
      by (simp add: mult-le-0-iff)
  qed
then show False
  using not-contained-13-2

```

```

    by simp
next
case neg-pos
have  $v *_R a \cdot h3 = a \cdot h1 + (-u) *_R a \cdot h2$ 
  using a1-eq
  by (simp add: algebra-simps)
have  $\{x. a \cdot h1 \cdot x \leq 0\} \cap \{x. a \cdot h2 \cdot x \leq 0\} \subseteq \{x. a \cdot h3 \cdot x \leq 0\}$ 
proof
  fix x
  assume x-in:  $x \in \{x. a \cdot h1 \cdot x \leq 0\} \cap \{x. a \cdot h2 \cdot x \leq 0\}$ 
  have  $v * (a \cdot h3 \cdot x) = (a \cdot h1 \cdot x) + (-u) * (a \cdot h2 \cdot x)$ 
    using  $\langle v *_R a \cdot h3 = a \cdot h1 + (-u) *_R a \cdot h2 \rangle$ 
    by (metis inner-left-distrib inner-scaleR-left)
  also have  $\dots \leq 0$ 
    using x-in neg-pos
    by (smt (verit, best) IntE mem-Collect-eq mult-le-0-iff)
  finally show  $x \in \{x. a \cdot h3 \cdot x \leq 0\}$ 
    using  $\langle v > 0 \rangle$ 
    by (simp add: mult-le-0-iff)
qed
then show False
  using not-contained-12-3
  by simp
next
case neg-neg
have  $a \cdot h1 \cdot w = u * (a \cdot h2 \cdot w) + v * (a \cdot h3 \cdot w)$ 
  using a1-eq
  by (simp add: inner-add)
moreover have  $u * (a \cdot h2 \cdot w) > 0 \wedge v * (a \cdot h3 \cdot w) > 0$ 
  using neg-neg w-props(2,3) mult-neg-neg
  by blast+
ultimately have  $a \cdot h1 \cdot w > 0$ 
  by simp
then show False
  using w-props(1) by simp
qed
qed

then have  $\dim (\text{span } \{a \cdot h1, a \cdot h2, a \cdot h3\}) = 3$ 
  using a-hs-neg(1,2,3) dim-span-eq-card-independent
  by fastforce
then show ?thesis
  using  $\langle \dim (\text{span } \{a \cdot h1, a \cdot h2, a \cdot h3\}) \leq 2 \rangle$ 
  by linarith
qed
qed

lemma polyhedron-ridge-two-facets:
  fixes p :: 'a :: euclidean-space set

```

assumes *polyhedron* p
and r *face-of* p
and $r \neq \{\}$
and $\text{aff-dim } r = \text{aff-dim } p - 2$
shows $\exists f1\ f2. f1 \text{ face-of } p \wedge \text{aff-dim } f1 = \text{aff-dim } p - 1 \wedge$
 $f2 \text{ face-of } p \wedge \text{aff-dim } f2 = \text{aff-dim } p - 1 \wedge$
 $f1 \neq f2 \wedge r \subseteq f1 \wedge r \subseteq f2 \wedge r = f1 \cap f2 \wedge$
 $(\forall f. f \text{ face-of } p \wedge \text{aff-dim } f = \text{aff-dim } p - 1 \wedge r \subseteq f$
 $\longrightarrow f = f1 \vee f = f2)$

proof –

obtain z **where** $z \in r$
using *assms*(3) **by** *blast*
then have $z \in p$
using *assms*(2) *face-of-imp-subset* **by** *auto*
let $?r' = ((+) (-z)) \text{ ' } r$
let $?p' = ((+) (-z)) \text{ ' } p$
have *origin-r'*: $0 \in ?r'$
using $\langle z \in r \rangle$ **by** *simp*
have *origin-p'*: $0 \in ?p'$
using $\langle z \in p \rangle$ **by** *simp*
have *polyhedron*: *polyhedron* $?p'$
using *assms*(1) *polyhedron-translation-eq*[*of* $p - z$] **by** *blast*
have *face-of*: $?r' \text{ face-of } ?p'$
using *assms*(2) *face-of-translation-eq* **by** *blast*
have *nonempty*: $?r' \neq \{\}$
using *assms*(3) **by** *blast*
have *aff-dims*: $\text{aff-dim } ?r' = \text{aff-dim } ?p' - 2$
by (*metis* *aff-dim-translation-eq* *assms*(4))
obtain $f1\ f2$ **where**
 $f1\text{-face-of}$: $f1 \text{ face-of } ?p'$ **and**
 aff-dim-f1 : $\text{aff-dim } f1 = \text{aff-dim } ?p' - 1$ **and**
 $f2\text{-face-of}$: $f2 \text{ face-of } ?p'$ **and**
 aff-dim-f2 : $\text{aff-dim } f2 = \text{aff-dim } ?p' - 1$ **and**
 $fs\text{-neg}$: $f1 \neq f2$ **and**
 $sub\text{-f1}$: $?r' \subseteq f1$ **and**
 $sub\text{-f2}$: $?r' \subseteq f2$ **and**
 $fs\text{-inter-eq}$: $f1 \cap f2 = ?r'$ **and**
 $only\text{-fs}$: $(\forall f. f \text{ face-of } ?p' \wedge$
 $\text{aff-dim } f = \text{aff-dim } ?p' - 1 \wedge ?r' \subseteq f \longrightarrow$
 $f = f1 \vee f = f2)$
using *polyhedron-ridge-two-facets-0*[*OF* *polyhedron face-of origin-r' aff-dims*]
by *auto*
let $?f1' = ((+) z) \text{ ' } f1$
let $?f2' = ((+) z) \text{ ' } f2$
have $?f1' \text{ face-of } p$
by (*metis* *f1-face-of face-of-translation-eq translation-galois*)
moreover have $\text{aff-dim } ?f1' = \text{aff-dim } p - 1$
by (*metis* *aff-dim-f1 aff-dim-translation-eq*)
moreover have $?f2' \text{ face-of } p$

by (metis f2-face-of face-of-translation-eq translation-galois)
 moreover have $\text{aff-dim } ?f2' = \text{aff-dim } p - 1$
 by (metis aff-dim-f2 aff-dim-translation-eq)
 moreover have $?f1' \neq ?f2'$
 using fs-neq translation-invert
 by auto
 moreover have $r \subseteq ?f1'$
 by (metis sub-f1 subset-image-iff translation-galois)
 moreover have $r \subseteq ?f2'$
 by (metis sub-f2 subset-image-iff translation-galois)
 moreover have $?f1' \cap ?f2' = r$
 by (metis fs-inter-eq translation-Int translation-galois)
 moreover have $(\forall f. f \text{ face-of } p \wedge$
 $\text{aff-dim } f = \text{aff-dim } p - 1 \wedge r \subseteq f \longrightarrow$
 $f = ?f1' \vee f = ?f2')$
 by (smt (verit, best) aff-dim-translation-eq face-of-translation-eq image-mono
 only-fs translation-galois)
 ultimately show ?thesis
 by blast
 qed

3.4 Lower bounds on face counts

3.4.1 Vertices

lemma *polytope-vertex-lower-bound*:

assumes *polytope* p
 shows $\text{aff-dim } p + 1 \leq \text{card } \{v. v \text{ face-of } p \wedge \text{aff-dim } v = 0\}$
 proof –
 let $?V = \{v. v \text{ extreme-point-of } p\}$
 have sets-eq: $\{v. v \text{ face-of } p \wedge \text{aff-dim } v = 0\} = (\lambda v. \{v\}) \text{ ` } ?V$
 using face-of-singleton aff-dim-sing aff-dim-eq-0
 by auto
 then have cards-eq: $\text{card } \{v. v \text{ face-of } p \wedge \text{aff-dim } v = 0\} = \text{card } ?V$
 unfolding sets-eq using card-image inj-singleton
 by blast
 have $\text{aff-dim } p + 1 \leq \text{aff-dim } (\text{convex hull } ?V) + 1$
 using Krein-Milman-Minkowski[OF *polytope-imp-compact* [OF *assms*] *polytope-imp-convex* [OF
assms]]
 by simp
 moreover have $\text{aff-dim } (\text{convex hull } ?V) + 1 \leq \text{card } ?V$
 using finite-polyhedron-extreme-points[OF *polytope-imp-polyhedron* [OF *assms*]]
 aff-dim-le-card[of ?V] aff-dim-convex-hull[of ?V]
 by simp
 ultimately show ?thesis
 unfolding cards-eq by auto
 qed

3.4.2 Facets

lemma *polytope-facet-lower-bound-0*:

assumes *polytope* p
assumes *aff-dim* $p \neq 0$
assumes *origin-in-p*: $0 \in p$
shows *aff-dim* $p + 1 \leq \text{card } \{f. f \text{ facet-of } p\}$

proof –

have *aff-dim-eq-dim*: *aff-dim* $p = \text{dim } p$
using *origin-in-p* *aff-dim-zero*[*of* p]
by (*simp* *add*: *hull-inc*)

obtain H a b **where**

finite H **and**
p-eq: $p = \text{affine hull } p \cap \bigcap H$ **and**
halfspaces: $\bigwedge h. h \in H \implies a \cdot h \neq 0 \wedge h = \{x. a \cdot h \cdot x \leq b \cdot h\} \wedge$
 $(\forall x \in \text{affine hull } p. x + a \cdot h \in \text{affine hull } p)$ **and**
minimal: $\bigwedge F'. F' \subset H \implies p \subset \text{affine hull } p \cap \bigcap F'$
using *polyhedron-Int-affine-parallel-minimal*[*of* p]
polytope-imp-polyhedron[*OF* *assms*(1)]
by *metis*

have *p-eq-span*: $p = \text{span } p \cap \bigcap H$
using *p-eq* $\langle 0 \in p \rangle$ *affine-hull-span-0*
by *blast*

have $\bigwedge h. h \in H \implies 0 \in h$
using $\langle 0 \in p \rangle$ *p-eq*
by *blast*

then have *b-nonneg*: $\bigwedge h. h \in H \implies 0 \leq b \cdot h$
by (*metis* *halfspaces* *inner-zero-right* *mem-Collect-eq*)

have *ineq1*: $\text{dim } p + 1 \leq \text{card } H$

proof (*rule ccontr*)

assume $\neg(\text{dim } p + 1 \leq \text{card } H)$
then have *: $\text{card } H \leq \text{dim } p$
by *auto*

have $H \neq \{\}$

by (*smt* (*verit*, *ccfv-threshold*) *IntE* *Int-lower2* *affine-affine-hull*
affine-bounded-eq-lowdim *closure-Inter-convex-open*
convex-rel-interior-finite-Inter *emptyE* *finite.intros*(1) *image-is-empty*
inf.orderE *local.aff-dim-eq-dim* *of-nat-0-le-iff* *origin-in-p* *p-eq*
polytope-eq-bounded-polyhedron *rel-interior-eq-closure* *subset-hull*
assms(1,2))

then obtain h **where** $h \in H$

by *blast*

have $\forall h' \in H. a \cdot h' \in \text{span } p$

by (*metis* *add-0* *affine-hull-span-0* *halfspaces* *hull-inc* *origin-in-p*)

then have $\text{span } (a \cdot (H - \{h\})) \subseteq \text{span } p$

using *span-mono*[*of* $a \cdot (H - \{h\})$ *span* p] *span-span*[*of* p]

by *blast*

moreover have $\text{span } (a \cdot (H - \{h\})) \neq \text{span } p$

```

proof
  assume  $\text{span } (a \cdot (H - \{h\})) = \text{span } p$ 
  then have  $\dim (\text{span } p) = \dim (\text{span } (a \cdot (H - \{h\})))$ 
    by arg0
  also have  $\dots \leq \text{card } (a \cdot (H - \{h\}))$ 
    by (simp add: <finite H> dim-le-card')
  finally show False
    by (metis (no-types, lifting) * <finite H> <h ∈ H> card-Diff1-less-iff
      card-image-le dim-span finite-Diff linorder-not-le order-trans)
qed
ultimately have  $\text{span-psubset: span } (a \cdot (H - \{h\})) \subset \text{span } p$ 
  by blast
obtain  $n$  where
   $n \neq 0$ 
   $n \in \text{span } p$  and
   $\text{orthogonal: } \bigwedge x. x \in \text{span } (a \cdot (H - \{h\})) \implies \text{orthogonal } n \ x$ 
  using orthogonal-to-subspace-exists-gen[OF span-psubset]
  by blast
obtain  $B$  where  $\text{bounded: } \bigwedge x. x \in p \implies \text{norm } x \leq B$ 
  using polytope-imp-bounded[OF assms(1)] bounded-pos[of p]
  by blast
{assume  $a \cdot h \cdot n \geq 0$ 
  let  $?x = -((B + 1) / (\text{norm } n)) * n$ 
  have  $?x \in \text{span } p$ 
    using  $\langle n \in \text{span } p \rangle$  span-mul
    by blast
  moreover have  $\bigwedge k. k \in H \implies (a \cdot k) \cdot ?x \leq (b \cdot k)$ 
  proof -
    fix  $k$ 
    assume  $k \in H$ 
    {assume  $k = h$ 
      have  $a \cdot h \cdot ?x \leq b \cdot h$ 
        using  $\langle 0 \leq a \cdot h \cdot n \rangle \langle h \in H \rangle$  b-nonneg bounded divide-nonneg-nonneg
        inner-simps(6) origin-in-p real-0-le-add-iff
        by fastforce
    } moreover {assume  $k \neq h$ 
      have  $(a \cdot k) \cdot ?x \leq (b \cdot k)$ 
        by (smt (verit, best) <k ∈ H> <k ≠ h> b-nonneg image-eqI inner-commute
          inner-simps(6) insert-Diff-single insert-absorb insert-iff mult-eq-0-iff
          orthogonal orthogonal-def span-superset subsetD)
    } ultimately show  $a \cdot k \cdot ?x \leq b \cdot k$ 
      by blast
  }
qed
ultimately have  $?x \in p$ 
  using p-eq-span halfspaces
  by blast
have  $\text{norm } ?x > B$ 
  by (simp add: <n ≠ 0>)
then have False

```

```

    using bounded[OF ‹?x ∈ p›]
    by linarith
  } moreover {assume a h · n ≤ 0
    let ?x = ((B + 1) / (norm n)) *R n
    have ?x ∈ span p
      using ‹n ∈ span p› span-mul
      by blast
    moreover have  $\bigwedge k. k \in H \implies a \cdot k \cdot ?x \leq b \cdot k$ 
    proof -
      fix k
      assume k ∈ H
      {assume k = h
        have a h · ?x ≤ b h
        by (smt (verit, best) ‹a h · n ≤ 0› ‹h ∈ H› b-nonneg bounded inner-simps(6)
          norm-ge-zero norm-zero origin-in-p real-scaleR-def split-scaleR-neg-le
          zero-le-divide-iff zero-less-one-class.zero-le-one)
      } moreover {assume k ≠ h
        have (a k) · ?x ≤ (b k)
        by (smt (verit, best) ‹k ∈ H› ‹k ≠ h› b-nonneg image-eqI inner-commute
          inner-simps(6) insert-Diff-single insert-absorb insert-iff mult-eq-0-iff
          orthogonal orthogonal-def span-superset subsetD)
      } ultimately show a k · ?x ≤ b k
      by blast
    qed
    ultimately have ?x ∈ p
      using p-eq-span halfspaces
      by blast
    have norm ?x > B
      by (simp add: ‹n ≠ 0›)
    then have False
      using bounded[OF ‹?x ∈ p›]
      by linarith
  } ultimately show False
    by linarith
qed

have ineq2: card H = card {f. f facet-of p}
proof -
  let ?g = λh. p ∩ {x. a h · x = b h}
  have halfspaces-simple:  $\bigwedge h'. h' \in H \implies a \cdot h' \neq 0 \wedge h' = \{x. a \cdot h' \cdot x \leq b \cdot h'\}$ 
    using halfspaces by blast
  have {f. f facet-of p} = ?g ‘ H
    using facet-of-polyhedron-explicit[OF ‹finite H› p-eq halfspaces-simple] mini-
mal
    by auto
  moreover have inj-on ?g H
    using facets-of-polyhedron-explicit-distinct[OF ‹finite H› p-eq halfspaces-simple]
    by (auto simp: minimal inj-on-def)
  ultimately show ?thesis

```

```

    using card-image by fastforce
qed

show ?thesis
  unfolding aff-dim-eq-dim
  using ineq1 ineq2
  by linarith
qed

lemma polytope-facet-lower-bound:
  assumes polytope p
  assumes aff-dim p  $\neq$  0
  shows aff-dim p + 1  $\leq$  card {f. f facet-of p}
proof (cases p = {})
  case True
  then show ?thesis
    using aff-dim-empty[of p] facet-of-empty
    by simp
next
  case False
  then obtain z where z  $\in$  p
    by blast
  have has-origin: 0  $\in$  ((+) (-z) ' p)
    using 'z  $\in$  p' by simp
  have is-polytope: polytope ((+) (-z) ' p)
    using polytope-translation-eq[of -z p] assms(1)
    by blast
  have nonzero-aff-dim: aff-dim ((+) (-z) ' p)  $\neq$  0
    using aff-dim-translation-eq[of -z p] assms(2)
    by argo
  have h1: card {f. f facet-of (+) (-z) ' p} = card ((') ((+) (-z)) ' {f. f facet-of
p})
    using facets-of-translation[of -z p]
    by argo
  moreover have inj-on ((') ((+) (-z))) {f. f facet-of p}
    by (simp add: inj-on-image)
  ultimately have card {f. f facet-of (+) (-z) ' p} = card {f. f facet-of p}
    by (metis card-image)
  then show ?thesis
    using polytope-facet-lower-bound-0[OF is-polytope nonzero-aff-dim has-origin]
    unfolding aff-dim-translation-eq[of -z p]
    by argo
qed

```

3.5 Double counting

```

lemma multiple-counting:
  fixes R :: 'a  $\Rightarrow$  'b  $\Rightarrow$  bool
  assumes finite s

```

```

assumes finite t
assumes  $\bigwedge x. x \in s \implies \text{card } \{y \in t. R \ x \ y\} = m$ 
assumes  $\bigwedge y. y \in t \implies \text{card } \{x \in s. R \ x \ y\} = n$ 
shows  $m * \text{card } s = n * \text{card } t$ 
proof –
  have lhs:  $(\sum x \in s. \sum y \in \{y \in t. R \ x \ y\}. 1) = m * \text{card } s$ 
    using assms(1,3) by simp
  have rhs:  $(\sum y \in t. \sum x \in \{x \in s. R \ x \ y\}. 1) = n * \text{card } t$ 
    using assms(2,4) by simp
  show ?thesis
    using sum.swap-restrict[OF assms(1) assms(2), of  $\lambda x \ y. (1::\text{nat}) \ R$ ]
    unfolding lhs rhs .
qed

```

3.5.1 $mF = 2E$

Let F be the number of facets and E be the number of ridges. If there are m ridges per facet, then $m * F = 2 * E$.

lemma *facet-ridge-relation*:

```

assumes polytope p
  and  $\text{aff-dim } p \geq 2$ 
  and edges-per-face:  $\forall f. f \text{ face-of } p \wedge \text{aff-dim } f = \text{aff-dim } p - 1 \implies$ 
     $\text{card } \{e. e \text{ face-of } p \wedge \text{aff-dim } e = \text{aff-dim } p - 2 \wedge e \subseteq f\} = m$ 
shows  $m * \text{card } \{f. f \text{ face-of } p \wedge \text{aff-dim } f = \text{aff-dim } p - 1\} =$ 
   $2 * \text{card } \{e. e \text{ face-of } p \wedge \text{aff-dim } e = \text{aff-dim } p - 2\}$ 
proof –
  let ?faces-set =  $\{f. f \text{ face-of } p \wedge \text{aff-dim } f = \text{aff-dim } p - 1\}$ 
  let ?edges-set =  $\{e. e \text{ face-of } p \wedge \text{aff-dim } e = \text{aff-dim } p - 2\}$ 
  have finite-faces: finite ?faces-set
    using finite-polytope-faces[OF  $\langle \text{polytope } p \rangle$ ]
    by fast
  have finite-edges: finite ?edges-set
    using finite-polytope-faces[OF  $\langle \text{polytope } p \rangle$ ]
    by fast
  have lhs-m:  $\bigwedge f. f \in ?\text{faces-set} \implies \text{card } \{e \in ?\text{edges-set}. e \subseteq f\} = m$ 
    using edges-per-face by auto
  have rhs-2:  $\bigwedge e. e \in ?\text{edges-set} \implies \text{card } \{f \in ?\text{faces-set}. e \subseteq f\} = 2$ 
proof –
  fix e
  assume *:  $e \in ?\text{edges-set}$ 
  have e face-of p
    using *
    by blast
  have aff-dim-e:  $\text{aff-dim } e = \text{aff-dim } p - 2$ 
    using * by auto
  then have  $e \neq \{\}$ 
    using assms(2) by auto
  have polyhedron p
    using polytope-imp-polyhedron[OF  $\langle \text{polytope } p \rangle$ ] .

```

```

obtain  $f1\ f2$  where
   $f1$ -props:  $f1$  face-of  $p$   $\text{aff-dim } f1 = \text{aff-dim } p - 1$   $e \subseteq f1$  and
   $f2$ -props:  $f2$  face-of  $p$   $\text{aff-dim } f2 = \text{aff-dim } p - 1$   $e \subseteq f2$  and
  distinct:  $f1 \neq f2$  and
  characterization:  $\forall f. f$  face-of  $p \wedge \text{aff-dim } f = \text{aff-dim } p - 1 \wedge e \subseteq f \longrightarrow f$ 
     $= f1 \vee f = f2$ 
  using polyhedron-ridge-two-facets[ $OF \langle \text{polyhedron } p \rangle \langle e \text{ face-of } p \rangle \langle e \neq \{\} \rangle$ 
     $\text{aff-dim-}e]$ 
  by auto
  have  $\{f \in ?\text{faces-set}. e \subseteq f\} = \{f1, f2\}$ 
  using characterization  $f1$ -props  $f2$ -props
  by auto
  then show  $\text{card } \{f \in ?\text{faces-set}. e \subseteq f\} = 2$ 
  using distinct
  by simp
qed
show ?thesis
  using multiple-counting[ $OF$  finite-faces finite-edges lhs-m rhs-2] .
qed

```

3.5.2 $nV = 2E$

Let V be the number of vertices and E be the number of edges. If there are n edges per vertex, then $n*V = 2*E$.

lemma edge-vertex-relation:

```

assumes polytope  $p$ 
  and edges-per-vert:  $\forall v. v$  face-of  $p \wedge \text{aff-dim } v = 0 \longrightarrow$ 
     $\text{card } \{e. e \text{ face-of } p \wedge \text{aff-dim } e = 1 \wedge v \subseteq e\} = n$ 
shows  $n * \text{card } \{v. v \text{ face-of } p \wedge \text{aff-dim } v = 0\} =$ 
   $2 * \text{card } \{e. e \text{ face-of } p \wedge \text{aff-dim } e = 1\}$ 
proof –
  let ?vertices-set =  $\{v. v \text{ face-of } p \wedge \text{aff-dim } v = 0\}$ 
  let ?edges-set =  $\{e. e \text{ face-of } p \wedge \text{aff-dim } e = 1\}$ 
  have finite-vertices: finite ?vertices-set
    using finite-polytope-faces[ $OF \langle \text{polytope } p \rangle$ ]
    by fast
  have finite-edges: finite ?edges-set
    using finite-polytope-faces[ $OF \langle \text{polytope } p \rangle$ ]
    by fast
  have lhs-n:  $\bigwedge v. v \in ?\text{vertices-set} \implies \text{card } \{e \in ?\text{edges-set}. v \subseteq e\} = n$ 
    using edges-per-vert by auto
  have rhs-2:  $\bigwedge e. e \in ?\text{edges-set} \implies \text{card } \{v \in ?\text{vertices-set}. v \subseteq e\} = 2$ 
proof –
  fix  $e$ 
  assume *:  $e \in ?\text{edges-set}$ 
  have  $\text{aff-dim } e = 1$ 
    using * by auto
  then have  $e \neq \{\}$ 
    by force

```

```

have compact e
  using * face-of-polytope-polytope[OF ‹polytope p›] polytope-imp-compact[of e]
  by blast
have convex e
  using * unfolding face-of-def
  by blast
have collinear e
proof -
  obtain B where
    aff-hull-B-e: affine hull B = affine hull e and
    ¬ affine-dependent B and
    int (card B) = 2
  using ‹aff-dim e = 1› aff-dim-basis-exists[of e] by auto
  then obtain x y where B-eq: B = {x, y} x ≠ y
  by (meson card-2-iff of-nat-eq-numeral-iff)
  have e ⊆ affine hull B
  by (simp add: aff-hull-B-e hull-subset)
  then show ?thesis
  unfolding B-eq using collinear-affine-hull
  by fast
qed
then obtain a b where e-eq: e = closed-segment a b
  using compact-convex-collinear-segment[OF ‹e ≠ {}› ‹compact e› ‹convex e›
‹collinear e›]
  by blast
then have a ≠ b
  using ‹aff-dim e = 1›
  by force

have {v. v extreme-point-of e} = {v. {v} face-of e ∧ aff-dim {v} = 0}
  unfolding face-of-singleton[symmetric] using aff-dim-sing
  by blast
then have (λx. {x}) ‘ {v. v extreme-point-of e} = {v. v face-of e ∧ aff-dim v
= 0}
  using aff-dim-eq-0 by auto
then have (λx. {x}) ‘ {v. v extreme-point-of e} = {v. v face-of p ∧ aff-dim v
= 0 ∧ v ⊆ e}
  using face-of-face * by force
then have bij-betw (λx. {x}) {v. v extreme-point-of e} {v ∈ ?vertices-set. v ⊆
e}
  by (simp add: bij-betw-imageI inj-on-def)
then have card {v ∈ ?vertices-set. v ⊆ e} = card {v. v extreme-point-of e}
  using bij-betw-same-card by force
moreover have {v. v extreme-point-of e} = {a, b}
  unfolding e-eq using extreme-point-of-segment[of - a b] ‹a ≠ b›
  by auto
moreover have card {a, b} = 2
  using ‹a ≠ b› by simp
ultimately show card {v ∈ ?vertices-set. v ⊆ e} = 2

```

by *argo*
 qed
 show *?thesis*
 using *multiple-counting[OF finite-vertices finite-edges lhs-n rhs-2]* .
 qed

hide-fact *multiple-counting*

3.6 Face containment

lemma *ridge-belongs-to-facet*:

assumes *e face-of p*
 and $\text{aff-dim } e = \text{aff-dim } p - 2$
 and $\text{aff-dim } p \geq 2$
 and *polyhedron p*
 obtains *f* where *f face-of p* $\text{aff-dim } f = \text{aff-dim } p - 1$ *e face-of f*
 proof –
 have $e \neq \{\}$
 using *assms(2,3)* by *auto*
 moreover have $e \neq p$
 using *assms(2,3)* by *auto*
 ultimately obtain *f* where *f facet-of p* $e \subseteq f$
 using *assms(1,4)* *face-of-polyhedron-subset-facet*
 by *blast*
 then have *f face-of p* $\text{aff-dim } f = \text{aff-dim } p - 1$
 unfolding *facet-of-def* by *simp+*
 moreover have *e face-of f*
 using $\langle f \text{ face-of } p \rangle$ *assms(1)* $\langle e \subseteq f \rangle$ *face-of-face*
 by *blast*
 ultimately show $(\bigwedge f. f \text{ face-of } p \implies \text{aff-dim } f = \text{aff-dim } p - 1 \implies e \text{ face-of } f) \implies \text{thesis} \implies \text{thesis}$
 by *blast*
 qed

lemma *peak-belongs-to-ridge*:

assumes *v face-of p*
 and $\text{aff-dim } v = \text{aff-dim } p - 3$
 and $\text{aff-dim } p \geq 3$
 and *polyhedron p*
 obtains *e* where *e face-of p* $\text{aff-dim } e = \text{aff-dim } p - 2$ *v face-of e*
 proof –
 have $v \neq \{\}$
 using *assms(2,3)* by *auto*
 moreover have $v \neq p$
 using *assms(2,3)* by *auto*
 ultimately obtain *f* where *f facet-of p* $v \subseteq f$
 using *assms(1,4)* *face-of-polyhedron-subset-facet*
 by *blast*

```

then have  $f \text{ face-of } p \text{ aff-dim } f = \text{aff-dim } p - 1$ 
  unfolding facet-of-def assms(3)
  by simp+
moreover have  $v \text{ face-of } f$ 
  using  $\langle f \text{ face-of } p \rangle \text{ assms}(1) \langle v \subseteq f \rangle \text{ face-of-face}$ 
  by blast
obtain  $e$  where  $e \text{ face-of } f \text{ aff-dim } e = \text{aff-dim } p - 2 \text{ } v \text{ face-of } e$ 
  using ridge-belongs-to-facet[OF  $\langle v \text{ face-of } f \rangle$  - -
    face-of-polyhedron-polyhedron[OF assms(4)  $\langle f \text{ face-of } p \rangle$ ]]
    assms(2,3)  $\langle \text{aff-dim } f = \text{aff-dim } p - 1 \rangle$ 
  by auto
then show  $(\bigwedge e. e \text{ face-of } p \implies \text{aff-dim } e = \text{aff-dim } p - 2 \implies v \text{ face-of } e \implies$ 
thesis)  $\implies$  thesis
  using face-of-trans[OF  $\langle e \text{ face-of } f \rangle \langle f \text{ face-of } p \rangle$ ]
  by blast
qed

end

```

4 Computational helpers

theory *Computation*

imports

Convex-More

Determinants-More

HOL-Analysis.Cross3

HOL-Analysis.Polytope

HOL-Eisbach.Eisbach

Polytope-More

HOL-Examples.Sqrt

begin

4.1 Operations in $\mathbb{Q}(\sqrt{5})$

Numbers of the form $a + b\sqrt{5}$, where $a, b \in \mathbb{Q}$

abbreviation $\varphi :: \text{real}$

where $\varphi \equiv (1 + \text{sqrt } 5)/2$

lemma *real-rat5-mul*:

```

 $(a1 + b1 * \text{sqrt } 5) * (a2 + b2 * \text{sqrt } 5)$ 
 $= (a1 * a2 + 5 * b1 * b2) + (a1 * b2 + a2 * b1) * \text{sqrt } 5$ 
by (simp add: algebra-simps)

```

lemma *real-rat5-inv*:

assumes $a^2 \neq 5 * b^2$

shows $\text{inverse } (a + b * \text{sqrt } 5) = a / (a^2 - 5 * b^2) + -b / (a^2 - 5 * b^2)$

$* \text{sqrt } 5$

proof -

```

have (a + b * sqrt 5) * (a - b * sqrt 5) = a^2 - 5 * b^2
  by (simp add: algebra-simps power2-eq-square)
moreover have a^2 - 5 * b^2 ≠ 0
  using assms by simp
ultimately have inverse (a + b * sqrt 5) = (a - b * sqrt 5) / (a^2 - 5 * b^2)
  by (metis divide-inverse mult-1 mult-eq-0-iff
    nonzero-divide-mult-cancel-right)
then show ?thesis
  by argo
qed

```

lemma *real-rat5-div*:

```

assumes a2^2 ≠ 5 * b2^2
shows (a1 + b1 * sqrt 5) / (a2 + b2 * sqrt 5) =
  ((a1 * a2 - 5 * b1 * b2) / (a2^2 - 5 * b2^2)) +
  ((a2 * b1 - a1 * b2) / (a2^2 - 5 * b2^2)) * sqrt 5
unfolding divide-inverse[of (a1 + b1 * sqrt 5) (a2 + b2 * sqrt 5)]
  real-rat5-inv[OF assms]
using real-rat5-mul[of a1 b1 a2 / (a2^2 - 5 * b2^2) - b2 / (a2^2 - 5 * b2^2)]
by argo

```

lemma *real-rat5-le-sqrt-cases*:

```

x ≤ y * sqrt 5 ⟷
(x ≤ 0 ∧ 0 ≤ y) ∨
(0 ≤ x ∧ 0 ≤ y ∧ x^2 ≤ 5 * y^2) ∨
(x ≤ 0 ∧ y ≤ 0 ∧ 5 * y^2 ≤ x^2)

```

proof –

```

have sqrt5-pos: sqrt 5 > 0 by simp
have sqrt5-sq: (sqrt 5)^2 = 5 by simp

```

consider

```

(x-ge0-y-ge0) x ≥ 0 y ≥ 0 |
(x-ge0-y-le0) x ≥ 0 y ≤ 0 |
(x-le0-y-ge0) x ≤ 0 y ≥ 0 |
(x-le0-y-le0) x ≤ 0 y ≤ 0
by argo

```

then show *?thesis*

proof (*cases*)

case *x-ge0-y-ge0*

then show *?thesis*

using *power-mono-iff*[of *x y * sqrt 5 2*]

by (*auto simp add: power-mult-distrib sqrt5-sq mult-ac*)

next

case *x-ge0-y-le0*

then show *?thesis*

```

by (smt (verit) mult.left-commute mult-le-0-iff mult-zero-left
  not-numeral-le-zero power2-eq-square power2-less-eq-zero-iff
  real-sqrt-le-0-iff zero-le-mult-iff)

```

```

next
  case  $x \leq 0 \wedge y \geq 0$ 
  then show ?thesis
    by (smt (verit, del-insts) sqrt5-pos power-mult-distrib real-sqrt-mult
        real-sqrt-pow2-iff)
next
  case  $x \leq 0 \wedge y \leq 0$ 
  have  $x \leq y * \sqrt{5} \iff -x \geq -y * \sqrt{5}$ 
  by linarith
  also have  $\dots \iff (-x)^2 \geq (-y * \sqrt{5})^2$ 
  by (smt (verit) x-le0-y-le0 sqrt5-pos mult-nonneg-nonneg pos2 power-mono-iff)
  also have  $\dots \iff x^2 \geq 5 * y^2$ 
  by (simp add: power-mult-distrib sqrt5-sq mult-ac)
  finally show ?thesis using  $\langle x \leq 0 \rangle \langle y \leq 0 \rangle$  by auto
qed
qed

```

```

lemma real-rat5-le-comparison:
  ( $a1 + b1 * \sqrt{5} \leq a2 + b2 * \sqrt{5}$ )  $\iff$ 
  ( $a1 \leq a2 \wedge b1 \leq b2$ )  $\vee$ 
  ( $a2 \leq a1 \wedge b1 \leq b2 \wedge (a1 - a2)^2 \leq 5 * (b2 - b1)^2$ )  $\vee$ 
  ( $a1 \leq a2 \wedge b2 \leq b1 \wedge 5 * (b2 - b1)^2 \leq (a1 - a2)^2$ )
  using real-rat5-le-sqrt-cases[of  $a1 - a2$   $b2 - b1$ ]
  by argo

```

Wrapper definition to avoiding oversimplification when applying simp

```

definition rat5 :: rat  $\Rightarrow$  rat  $\Rightarrow$  real
  where rat5 a b = real-of-rat a + real-of-rat b * sqrt 5

```

```

lemma rat5-add: rat5 a1 b1 + rat5 a2 b2 = rat5 (a1 + a2) (b1 + b2)
  unfolding rat5-def
  by (simp add: algebra-simps of-rat-add)

```

```

lemma rat5-sub: rat5 a1 b1 - rat5 a2 b2 = rat5 (a1 - a2) (b1 - b2)
  unfolding rat5-def
  by (simp add: algebra-simps of-rat-diff)

```

```

lemma rat5-mul: rat5 a1 b1 * rat5 a2 b2 = rat5 (a1 * a2 + 5 * b1 * b2) (a1 *
  b2 + b1 * a2)
  unfolding rat5-def
  by (simp add: algebra-simps of-rat-mult of-rat-add)

```

```

lemma rat5-div:
  assumes  $a2 \neq 0 \vee b2 \neq 0$ 
  shows rat5 a1 b1 / rat5 a2 b2 =
    rat5 ((a1 * a2 - 5 * b1 * b2) / (a2^2 - 5 * b2^2))
      ((a2 * b1 - a1 * b2) / (a2^2 - 5 * b2^2))

```

```

proof -
  have denom-nonzero: (real-of-rat a2)^2  $\neq$  5 * (real-of-rat b2)^2

```

```

proof
  assume *: (real-of-rat a2)2 = 5 * (real-of-rat b2)2
  show False
  proof (cases b2 = 0)
    case True
    then show False
    using * assms by simp
  next
  case False
  have 5 = (real-of-rat a2 / real-of-rat b2)2
  using * False
  by (simp add: field-simps of-rat-divide)
  then have sqrt 5 = abs (real-of-rat (a2 / b2))
  by (metis of-rat-divide real-sqrt-abs)
  moreover have sqrt 5 ∉ ℚ
  using sqrt-prime-irrational[of 5] by auto
  moreover have abs (real-of-rat (a2 / b2)) ∈ ℚ
  by simp
  ultimately show False
  by auto
qed
qed
show ?thesis
  using rat5-def real-rat5-div[OF denom-nonzero]
  by (simp add: of-rat-diff of-rat-divide of-rat-mult of-rat-power)
qed

```

```

lemma rat5-le:
  (rat5 a1 b1 ≤ rat5 a2 b2) =
    ((a1 ≤ a2 ∧ b1 ≤ b2) ∨
     (a2 ≤ a1 ∧ b1 ≤ b2 ∧ (a1 - a2)2 ≤ 5 * (b2 - b1)2) ∨
     (a1 ≤ a2 ∧ b2 ≤ b1 ∧ 5 * (b2 - b1)2 ≤ (a1 - a2)2))
  unfolding rat5-def real-rat5-le-comparison
  by (metis (mono-tags, opaque-lifting) of-rat-diff of-rat-less-eq of-rat-mult of-rat-numeral-eq of-rat-power)

```

```

lemma rat5-eq: rat5 a1 b1 = rat5 a2 b2 ⟷ a1 = a2 ∧ b1 = b2

```

```

proof -
  {
    assume *: rat5 a1 b1 = rat5 a2 b2 b1 ≠ b2
    have real-of-rat (a1 - a2) = real-of-rat (b2 - b1) * sqrt 5
    using * rat5-def
    by (simp add: left-diff-distrib of-rat-diff)
    moreover have real-of-rat (b2 - b1) ≠ 0
    using * by simp
    ultimately have sqrt 5 = real-of-rat ((a1 - a2) / (b2 - b1))
    by (simp add: of-rat-divide)
    then have False
    using sqrt-prime-irrational[of 5] by auto
  }

```

```

} then show ?thesis
  using rat5-def by force
qed

```

```

lemma rat5-0-1:
  shows rat5 0 0 = 0 rat5 1 0 = 1 unfolding rat5-def by auto

```

```

lemma rat5-eq-0:
  rat5 a1 b1 = 0  $\longleftrightarrow$  a1 = 0  $\wedge$  b1 = 0
  using rat5-eq[of a1 b1 0 0] unfolding rat5-def
  by simp

```

```

lemma rat5-eq-1:
  rat5 a1 b1 = 1  $\longleftrightarrow$  a1 = 1  $\wedge$  b1 = 0
  using rat5-eq[of a1 b1 1 0] unfolding rat5-def
  by simp

```

4.2 Operations in $\mathbb{R}^3$

```

lemma vector3-add:
  fixes x1 x2 x3 y1 y2 y3 :: real
  shows vector [x1, x2, x3] + vector [y1, y2, y3] =
    vector [x1 + y1, x2 + y2, x3 + y3]
  by (simp add: vec-eq-iff vector-def forall-3)

```

```

lemma vector3-sub:
  fixes x1 x2 x3 y1 y2 y3 :: real
  shows vector [x1, x2, x3] - vector [y1, y2, y3] =
    vector [x1 - y1, x2 - y2, x3 - y3]
  by (simp add: vec-eq-iff vector-def forall-3)

```

unbundle cross3-syntax

```

lemma vector3-cross:
  fixes x1 x2 x3 y1 y2 y3 :: real
  shows vector [x1, x2, x3]  $\times$  vector [y1, y2, y3] =
    (vector [x2 * y3 - x3 * y2, x3 * y1 - x1 * y3, x1 * y2 - x2 * y1] ::
    real^3)
  by (simp add: cross3-def)

```

```

lemma vector3-eq-0:
  fixes x1 x2 x3 :: real
  shows (vector [x1, x2, x3] :: real^3) = 0  $\longleftrightarrow$  x1 = 0  $\wedge$  x2 = 0  $\wedge$  x3 = 0
  by (simp add: vec-eq-iff vector-def forall-3)

```

```

lemma vector3-eq:
  fixes x1 x2 x3 y1 y2 y3 :: real
  shows (vector [x1, x2, x3] :: real^3) = vector [y1, y2, y3]  $\longleftrightarrow$  x1 = y1  $\wedge$  x2
= y2  $\wedge$  x3 = y3

```

by (simp add: vec-eq-iff vector-def forall-3)

lemma vector3-neg:

shows $((\text{vector } [x1, x2, x3] :: \text{real}^3) \neq \text{vector } [y1, y2, y3]) =$
 $(x1 \neq y1 \vee x2 \neq y2 \vee x3 \neq y3)$
 by (simp add: vec-eq-iff forall-3)

lemma vector3-dot:

fixes $x1\ x2\ x3\ y1\ y2\ y3 :: \text{real}$
 shows $(\text{vector } [x1, x2, x3] :: \text{real}^3) \cdot \text{vector } [y1, y2, y3] =$
 $x1*y1 + x2*y2 + x3*y3$
 by (simp add: inner-vec-def sum-3 vector-def)

4.3 Computation of faces

4.3.1 2-faces

lemma compute-faces-2:

fixes $F\ S :: (\text{real}^3)\ \text{set}$
 assumes *finite S*
 shows $(F\ \text{face-of}\ (\text{convex hull } S) \wedge \text{aff-dim } F = 2) =$
 $(\exists x \in S. \exists y \in S. \exists z \in S.$
 $\text{let } a = (z - x) \times (y - x)\ \text{in}$
 $a \neq 0 \wedge$
 $(\text{let } b = a \cdot x\ \text{in}$
 $((\forall w \in S. a \cdot w \leq b) \vee (\forall w \in S. a \cdot w \geq b)) \wedge$
 $F = \text{convex hull } (S \cap \{x. a \cdot x = b\})))$
 (is ?lhs = ?rhs)

proof

assume ?lhs

then have *: $F\ \text{face-of}\ (\text{convex hull } S)\ \text{aff-dim } F = 2$ by blast+

then obtain T where $T\text{-facts: } T \subseteq S\ F = \text{convex hull } T$

by (meson assms face-of-convex-hull-subset finite-imp-compact)

then obtain U where $U\text{-facts: } U \subseteq T \neg \text{affine-dependent } U\ \text{affine hull } T =$
 $\text{affine hull } U$

using affine-basis-exists[of T] by blast

then have $\text{aff-dim } U = 2$

by (metis *(2) $T\text{-facts}(2)$ aff-dim-affine-hull aff-dim-convex-hull)

then have $\text{card } U = 3$

using $U\text{-facts}(2)$ aff-dim-affine-independent by fastforce

then obtain $x\ y\ z$ where

$U\text{-eq: } U = \{x, y, z\}$ and

$xyz\text{-neq: } x \neq y\ x \neq z\ y \neq z$

using card-3-iff[of U] by blast

then have $\text{not-collinear: } \neg \text{collinear } \{x, y, z\}$

using $U\text{-facts}(2)$ exposed-face-of-polyhedron

by (simp add: collinear-3-eq-affine-dependent)

have $xyz\text{-in-}S: x \in S\ y \in S\ z \in S$

using $U\text{-eq } U\text{-facts}(1)\ T\text{-facts}(1)$ by blast+

```

let ?a = (z - x) × (y - x)
let ?b = ?a · x
let ?H = {w. ?a · w = ?b}
have ?a ≠ 0
  using not-collinear collinear-3 cross-eq-0[symmetric]
  by (metis (no-types, opaque-lifting) NO-MATCH-def insert-commute)
have ?a · y = ?b
  by (metis (no-types, opaque-lifting) add.commute diff-0-right diff-add-cancel
dot-cross-self(2) inner-commute
inner-diff-left)
have ?a · z = ?b
  by (metis Cross3.left-diff-distrib cross-refl cross-triple diff-left-imp-eq diff-self)

have F exposed-face-of (convex hull S)
  using *(1) polyhedron-convex-hull[OF assms] exposed-face-of-polyhedron
  by blast
then obtain a' b' where a'-b'-facts:
  convex hull S ⊆ {w. a' · w ≤ b'}
  convex hull T = convex hull S ∩ {w. a' · w = b'}
  unfolding exposed-face-of-def T-facts(2) by blast
let ?H' = {w. a' · w = b'}

have {x, y, z} ⊆ ?H
  using ⟨?a · y = ?b⟩ ⟨?a · z = ?b⟩ by simp
then have affine hull {x, y, z} ⊆ ?H
  using affine-hyperplane hull-minimal by metis
then have affine hull T ⊆ ?H
  using U-eq U-facts(3) by blast
moreover have affine hull T ⊆ ?H'
  using a'-b'-facts(2)
  by (metis affine-hyperplane affine-imp-convex inf.cobounded2 subset-hull)
ultimately have affine hull T ⊆ ?H ∩ ?H'
  by blast
moreover have aff-dim (affine hull T) = 2
  by (metis U-facts(3) ⟨aff-dim U = 2⟩ aff-dim-affine-hull)
ultimately have aff-dim (?H ∩ ?H') ≥ 2
  using aff-dim-subset
  by metis

show ?rhs
proof (cases ?H' = UNIV)
  case True
    have S ⊆ affine hull T
      using True a'-b'-facts(2) convex-hull-subset-affine-hull hull-subset
      by fast
    then have s-in-plane: S ⊆ ?H
      using ⟨affine hull T ⊆ ?H⟩ by blast
    then have one-side: (∀ w ∈ S. ?a · w ≤ ?b) ∨ (∀ w ∈ S. ?a · w ≥ ?b)
      by auto

```

```

have  $S \cap ?H = S$ 
  using s-in-plane by auto
then have  $F = \text{convex hull } (S \cap \{w. ?a \cdot w = ?b\})$ 
  using True a'-b'-facts(2) T-facts(2) by simp

then show ?thesis
  using xyz-in-S  $\langle ?a \neq 0 \rangle$  one-side
  unfolding Let-def by blast
next
case False
have planes-eq:  $?H = ?H'$ 
proof (rule ccontr)
  assume neg:  $?H \neq ?H'$ 
  have  $\text{aff-dim } ?H = 2$ 
    using  $\langle ?a \neq 0 \rangle$  aff-dim-hyperplane[of  $?a ?b$ ] by simp
  then have  $\text{aff-dim } (?H \cap ?H') < 2$ 
    using aff-dim-affine-Int-hyperplane[of  $?H a' b'$ ]
      neg False subset-hyperplanes[of  $?a ?b a' b'$ ]
    by (simp add: affine-hyperplane)
  then show False
    using  $\langle \text{aff-dim } (?H \cap ?H') \geq 2 \rangle$ 
    by simp
qed

then have F-eq:  $F = \text{convex hull } (S \cap ?H)$ 
  by (smt (verit, ccfv-threshold) Int-commute T-facts(1,2) a'-b'-facts(2) con-
  vex-convex-hull hull-minimal hull-subset inf.orderE le-inf-iff subsetI)

have  $?H' \subseteq ?H$ 
  using planes-eq by auto
moreover have  $a' \neq 0$ 
  by (metis  $\langle ?a \neq 0 \rangle$  hyperplane-eq-UNIV hyperplane-eq-empty planes-eq)
ultimately obtain  $k$  where  $?a = k *_R a' ?b = k *_R b' k \neq 0$ 
  using  $\langle ?a \neq 0 \rangle$  hyperplane-subset-imp-scaled
  by blast
then have  $(\forall w \in S. ?a \cdot w \leq ?b) \vee (\forall w \in S. ?a \cdot w \geq ?b)$ 
  using a'-b'-facts(1) hull-subset[of  $S$ ]
  by (auto simp: mult-le-cancel-left)

then show ?thesis
  using xyz-in-S  $\langle ?a \neq 0 \rangle$  F-eq unfolding Let-def
  by blast
qed
next
assume ?rhs
then obtain  $x y z$  where
  xyz-in-S:  $x \in S y \in S z \in S$  and

```

$a\text{-neq-}0: (z - x) \times (y - x) \neq 0$ **and**
 $one\text{-}side: (\forall w \in S. (z - x) \times (y - x) \cdot w \leq (z - x) \times (y - x) \cdot x) \vee$
 $(\forall w \in S. (z - x) \times (y - x) \cdot w \geq (z - x) \times (y - x) \cdot x)$ **and**
 $F\text{-}eq: F = \text{convex hull } (S \cap \{w. (z - x) \times (y - x) \cdot w = (z - x) \times (y - x) \cdot$
 $x\})$
unfolding *Let-def* **by** *blast*
let $?a = (z - x) \times (y - x)$
let $?b = ?a \cdot x$
let $?H = \{w. ?a \cdot w = ?b\}$
have $\text{convex hull } (S \cap ?H) = \text{convex hull } S \cap ?H$
proof
show $\text{convex hull } (S \cap ?H) \subseteq \text{convex hull } S \cap ?H$
using *convex-hyperplane*[*of* $?a$ $?b$] *hull-mono*[*of* $S \cap ?H$ S *convex*]
subset-hull[*of* $\text{convex } ?H$ $S \cap ?H$]
by *blast*
next
show $\text{convex hull } S \cap ?H \subseteq \text{convex hull } (S \cap ?H)$
proof (*cases* $S \subseteq ?H$)
case *True*
then show $?thesis$ **by** (*simp add: Int-absorb2*)
next
case *False*
let $?S\text{-}on = S \cap ?H$
let $?S\text{-}off = S - ?H$
have $\text{convex hull } S \cap ?H = \text{convex hull } (\text{convex hull } (?S\text{-}on) \cup \text{convex hull } (?S\text{-}off)) \cap ?H$
by (*metis Int-Diff-Un hull-Un-left hull-Un-right*)
moreover have $\text{convex hull } (\text{convex hull } (?S\text{-}on) \cup \text{convex hull } (?S\text{-}off)) \cap ?H \subseteq \text{convex hull } (S \cap ?H)$
proof
fix w
assume $*$: $w \in \text{convex hull } (\text{convex hull } (?S\text{-}on) \cup \text{convex hull } (?S\text{-}off)) \cap ?H$
have $\text{convex: convex } (\text{convex hull } (?S\text{-}on))$ $\text{convex } (\text{convex hull } (?S\text{-}off))$
using *convex-convex-hull* **by** *blast+*
have $\text{nonempty1: convex hull } ?S\text{-}on \neq \{\}$
using *xyz-in-S(1) convex-hull-eq-empty* **by** *blast*
have $\text{nonempty2: convex hull } ?S\text{-}off \neq \{\}$
using *False convex-hull-eq-empty* **by** *simp*
obtain u p q **where**
 $w\text{-facts: } w = (1 - u) *_R p + u *_R q$ $?a \cdot w = ?b$ **and**
 $p\text{-in: } p \in \text{convex hull } ?S\text{-}on$ **and**
 $q\text{-in: } q \in \text{convex hull } ?S\text{-}off$ **and**
 $u\text{-range: } 0 \leq u$ $u \leq 1$
using $*$ *convex-hull-union-nonempty-explicit*[*OF* *convex(1)* *nonempty1* *convex(2)* *nonempty2*]
by *blast*
show $w \in \text{convex hull } (S \cap ?H)$
proof (*cases* $u = 0$)

```

case True
then show ?thesis
  using p-in w-facts(1) by simp
next
case u-nonzero: False
have  $p \in \text{convex hull } ?H$ 
  using p-in hull-mono[of ?S-on ?H] by blast
then have  $?a \cdot p = ?b$ 
  using convex-hyperplane hull-same
  by blast
have  $q \notin ?H$ 
proof
  assume  $q \in ?H$ 
  then have  $?a \cdot q = ?b$  by simp
  consider
    (leq)  $\forall w \in S. ?a \cdot w \leq ?b$  |
    (geq)  $\forall w \in S. ?a \cdot w \geq ?b$ 
  using one-side by blast
  then show False
  proof cases
    case leq
    then have  $?S\text{-off} \subseteq \{w. ?a \cdot w < ?b\}$ 
    by fastforce
    then have  $\text{convex hull } ?S\text{-off} \subseteq \{w. ?a \cdot w < ?b\}$ 
    using convex-halfspace-lt hull-minimal by metis
    then show False
    using  $\langle ?a \cdot q = ?b \rangle \langle q \in \text{convex hull } ?S\text{-off} \rangle$  by auto
  next
    case geq
    then have  $?S\text{-off} \subseteq \{w. ?a \cdot w > ?b\}$ 
    by fastforce
    then have  $\text{convex hull } ?S\text{-off} \subseteq \{w. ?a \cdot w > ?b\}$ 
    using convex-halfspace-gt hull-minimal by metis
    then show False
    using  $\langle ?a \cdot q = ?b \rangle \langle q \in \text{convex hull } ?S\text{-off} \rangle$  by auto
  qed
qed
then have  $?a \cdot q \neq ?b$  by simp
moreover have  $?a \cdot ((1 - u) *_{\mathbb{R}} p + u *_{\mathbb{R}} q) = (1 - u) * (?a \cdot p) + u$ 
 $* (?a \cdot q)$ 
  by (simp add: inner-add-right)
ultimately show ?thesis
  using  $\langle q \in \text{convex hull } ?S\text{-off} \rangle$ 
  by (metis  $\langle ?a \cdot p = ?b \rangle$  real-scaleR-def segment-degen-0 u-nonzero w-facts)
qed
qed
ultimately show ?thesis by argo
qed
qed

```

then have $F\text{-as-inter}$: $F = \text{convex hull } S \cap ?H$
using $F\text{-eq}$ **by** simp
consider $(\text{case-le}) \forall w \in S. ?a \cdot w \leq ?b \mid (\text{case-ge}) \forall w \in S. ?a \cdot w \geq ?b$
using one-side **by** blast
then have face-of : $F \text{ face-of } (\text{convex hull } S)$
proof cases
case case-le
then have subset : $S \subseteq \{x. ?a \cdot x \leq ?b\}$
by blast
have $\text{convex hull } S \subseteq \{x. ?a \cdot x \leq ?b\}$
using $\text{convex-halfspace-le}$ [of $?a ?b$] hull-minimal [OF subset , of convex]
by blast
then have $\bigwedge w. w \in \text{convex hull } S \implies ?a \cdot w \leq ?b$
by blast
then have $\text{convex hull } S \cap ?H \text{ face-of } \text{convex hull } S$
using $\text{face-of-Int-supporting-hyperplane-le}$ [OF $\text{convex-convex-hull}$ [of S]]
by blast
then show $?thesis$
using $F\text{-as-inter}$ **by** simp
next
case case-ge
then have subset : $S \subseteq \{x. ?a \cdot x \geq ?b\}$
by blast
have $\text{convex hull } S \subseteq \{x. ?a \cdot x \geq ?b\}$
using $\text{convex-halfspace-ge}$ [of $?b ?a$] hull-minimal [OF subset , of convex]
by blast
then have $\bigwedge w. w \in \text{convex hull } S \implies ?a \cdot w \geq ?b$
by blast
then have $\text{convex hull } S \cap ?H \text{ face-of } \text{convex hull } S$
using $\text{face-of-Int-supporting-hyperplane-ge}$ [OF $\text{convex-convex-hull}$ [of S]]
by blast
then show $?thesis$
using $F\text{-as-inter}$ **by** simp
qed

have $\text{aff-dim } ?H = 2$
using $\text{aff-dim-hyperplane}$ [OF $a\text{-neq-0}$, of $?b$] **by** simp
have $F \subseteq ?H$
using $F\text{-as-inter}$ **by** blast
then have upper : $\text{aff-dim } F \leq 2$
using aff-dim-subset [of $F ?H$] **unfolding** $\langle \text{aff-dim } ?H = 2 \rangle$
by blast

have $x \neq y \ x \neq z \ y \neq z$
using $a\text{-neq-0}$ **by** force+
then have $\text{card } \{x, y, z\} = 3$
by auto
have $\neg \text{collinear } \{x, y, z\}$
using collinear-3 [of $z \ x \ y$] cross-eq-0 $a\text{-neq-0}$

```

    by (simp add: insert-commute)
  then have dim-xyz: aff-dim {x, y, z} = 2
    using collinear-3-eq-affine-dependent aff-dim-affine-independent
    unfolding ⟨card {x, y, z} = 3⟩
    by fastforce
  have ?a · y = ?b
    by (metis (no-types, opaque-lifting) dot-cross-self(2) inner-commute inner-diff-left
right-minus-eq)
  then have y-in: y ∈ S ∩ ?H
    using xyz-in-S(2) by blast
  have ?a · z = ?b
    by (metis (no-types, lifting) Cross3.right-diff-distrib add-0 add-uminus-conv-diff
cross-refl cross-skew cross-triple diff-0-right)
  then have z-in: z ∈ S ∩ ?H
    using xyz-in-S(3) by blast
  then have {x, y, z} ⊆ convex hull (S ∩ ?H)
    using xyz-in-S(1) y-in hull-subset by fast
  then have xyz-in-F: {x, y, z} ⊆ F
    using F-eq by simp
  have lower: 2 ≤ aff-dim F
    using aff-dim-subset[OF xyz-in-F] dim-xyz by simp

```

```

show ?lhs
  using face-of upper lower by simp
qed

```

```

lemma set-choice-decomposition-1:
  assumes sym-yz:  $\bigwedge x y z. Q x y z \implies Q x z y$ 
  and sym-xy:  $\bigwedge x y z. Q x y z \implies Q y x z$ 
  and degen:  $\bigwedge x z. \neg Q x x z$ 
  shows  $(\exists x \in \text{insert } v S. \exists y \in \text{insert } v S. \exists z \in \text{insert } v S. Q x y z) \longleftrightarrow$ 
 $(\exists y \in S. \exists z \in S. Q v y z) \vee (\exists x \in S. \exists y \in S. \exists z \in S. Q x y z)$ 
  by (metis degen insertCI insertE sym-xy sym-yz)

```

```

lemma compute-faces-2-step-1:
  fixes F S v and T :: (real^3) set
  defines Q  $\equiv (\lambda x y z. \text{let } a = (z - x) \times (y - x) \text{ in}$ 
 $a \neq 0 \wedge$ 
 $(\text{let } b = a \cdot x \text{ in}$ 
 $((\forall w \in T. a \cdot w \leq b) \vee (\forall w \in T. a \cdot w \geq b)) \wedge$ 
 $F = \text{convex hull } (T \cap \{k. a \cdot k = b\})))$ 
  shows  $(\exists x \in (\text{insert } v S). \exists y \in (\text{insert } v S). \exists z \in (\text{insert } v S). Q x y z) =$ 
 $((\exists y \in S. \exists z \in S. Q v y z) \vee (\exists x \in S. \exists y \in S. \exists z \in S. Q x y z))$ 
  (is ?lhs = ?rhs)
proof -
  have Q-degen:  $\bigwedge x z. \neg Q x x z$ 
    unfolding Q-def Let-def
    by (simp add: cross3-simps)
  have Q-invariant:  $Q x y z \text{ if } Q x z y \text{ for } x y z$ 

```

using *that cross-skew*[of $y - x \ z - x$] **unfolding** *Q-def Let-def*
by *auto*

have *sym-xy*: $\bigwedge x \ y \ z. \ Q \ x \ y \ z \implies Q \ y \ x \ z$

proof $-$

fix $x \ y \ z$

assume $Q \ x \ y \ z$

then obtain $a \ b$ **where**

a-def: $a = (z - x) \times (y - x)$ **and**

a-nonzero: $a \neq 0$ **and**

b-def: $b = a \cdot x$ **and**

ineq: $(\forall w \in T. \ a \cdot w \leq b) \vee (\forall w \in T. \ a \cdot w \geq b)$ **and**

plane: $F = \text{convex hull } (T \cap \{w. \ a \cdot w = b\})$

unfolding *Q-def Let-def*

by *auto*

let $?a' = (z - y) \times (x - y)$

let $?b' = ?a' \cdot y$

have $?a' = -a$

by (*smt* (*verit*, *del-insts*) *Cross3.right-diff-distrib a-def cross-refl*
cross-skew eq-iff-diff-eq-0)

have *a-orth*: $a \cdot (y - x) = 0$

unfolding *a-def dot-cross-self*(3)

by *blast*

then have $a \cdot y = b$

unfolding *b-def*

by (*simp add: inner-diff-right*)

show $Q \ y \ x \ z$

unfolding *Q-def Let-def*

using $\langle a \cdot y = b \rangle \ \langle ?a' = -a \rangle \ a\text{-nonzero} \ \text{ineq} \ \text{plane}$

by *auto*

qed

show *?thesis*

using *set-choice-decomposition-1*[of $Q \ v \ S$] *Q-invariant sym-xy Q-degen Q-def*

by *blast*

qed

lemma *set-choice-decomposition-2*:

assumes *sym*: $\bigwedge x \ y. \ Q \ x \ y \implies Q \ y \ x$

and *degen*: $\bigwedge x. \ \neg \ Q \ x \ x$

shows $(\exists y \in \text{insert } u \ S. \ \exists z \in \text{insert } u \ S. \ Q \ y \ z) \longleftrightarrow$

$(\exists z \in S. \ Q \ u \ z) \vee (\exists y \in S. \ \exists z \in S. \ Q \ y \ z)$

by (*metis degen insertCI insertE sym*)

lemma *compute-faces-2-step-2*:

fixes $F \ u \ v \ S$ **and** $T :: (\text{real}^3) \ \text{set}$

defines $Q \equiv (\lambda y \ z. \ \text{let } a = (z - v) \times (y - v) \ \text{in}$

$a \neq 0 \wedge$

$(\text{let } b = a \cdot v \ \text{in}$

```

      (( $\forall w \in T. a \cdot w \leq b$ )  $\vee$  ( $\forall w \in T. a \cdot w \geq b$ ))  $\wedge$ 
       $F = \text{convex hull } (T \cap \{k. a \cdot k = b\})$ )
shows ( $\exists y \in (\text{insert } u \ S). \exists z \in (\text{insert } u \ S). Q \ y \ z =$ 
      ( $(\exists z \in S. Q \ u \ z) \vee (\exists y \in S. \exists z \in S. Q \ y \ z)$ )
(is ?lhs = ?rhs)
proof -
  have Q-degen:  $\bigwedge x. \neg Q \ x \ x$ 
    unfolding Q-def Let-def
    by (simp add: cross3-simps)

  have Q-sym:  $\bigwedge x \ y. Q \ x \ y \implies Q \ y \ x$ 
proof -
  fix x y
  assume Q x y
  then obtain a b where
    a-def:  $a = (y - v) \times (x - v)$  and
    a-nonzero:  $a \neq 0$  and
    b-def:  $b = a \cdot v$  and
    ineq: ( $\forall w \in T. a \cdot w \leq b$ )  $\vee$  ( $\forall w \in T. a \cdot w \geq b$ ) and
    plane:  $F = \text{convex hull } (T \cap \{w. a \cdot w = b\})$ 
    unfolding Q-def Let-def
    by auto
  let ?a' =  $(x - v) \times (y - v)$ 
  let ?b' =  $?a' \cdot v$ 
  have ?a' =  $-a$ 
    unfolding a-def using cross-skew
    by blast
  have  $\{w. ?a' \cdot w = ?b'\} = \{w. a \cdot w = b\}$ 
    unfolding  $\langle ?a' = -a \rangle$  using b-def
    by auto
  then show Q y x
    unfolding Q-def Let-def
    using  $\langle ?a' = -a \rangle$  a-nonzero plane ineq
    by auto
qed

show ?thesis
  using set-choice-decomposition-2[of Q u S] Q-sym Q-degen Q-def
  by blast
qed

```

4.3.2 1-faces

```

lemma compute-faces-1:
  fixes F n d and S :: ( $\text{real}^{\mathcal{Z}}$ ) set
  assumes  $\bigwedge x. x \in S \implies n \cdot x = d$ 
    and finite S
    and  $n \neq 0$ 
  shows ( $F \text{ face-of } (\text{convex hull } S) \wedge \text{aff-dim } F = 1$ ) =

```

```

    (∃ x∈S. ∃ y∈S.
      let a = n × (y - x) in
      a ≠ 0 ∧
      (let b = a · x in
        ((∀ w∈S. a · w ≤ b) ∨ (∀ w∈S. a · w ≥ b)) ∧
        F = convex hull (S ∩ {w. a · w = b})))
  (is ?lhs = ?rhs)
proof
  assume ?lhs
  then have *: F face-of (convex hull S) aff-dim F = 1 by blast+
  then obtain T where T-facts: T ⊆ S F = convex hull T
    by (meson assms face-of-convex-hull-subset finite-imp-compact)
  then obtain U where U-facts: U ⊆ T ¬ affine-dependent U affine hull T =
    affine hull U
    using affine-basis-exists[of T] by blast
  then have aff-dim U = 1
    by (metis *(2) T-facts(2) aff-dim-affine-hull aff-dim-convex-hull)
  then have card U = 2
    using U-facts(2) aff-dim-affine-independent by fastforce
  then obtain x y where U-eq: U = {x, y} and xy-neq: x ≠ y
    using card-2-iff[of U] by blast
  have xy-in-S: x ∈ S y ∈ S
    using U-eq U-facts(1) T-facts(1) by blast+

  let ?a = n × (y - x)
  let ?b = ?a · x
  let ?H = {w. ?a · w = ?b}

  have n · (y - x) = 0
    using assms(1)[OF xy-in-S(1)] assms(1)[OF xy-in-S(2)]
    by (simp add: inner-diff-right)
  then have ?a ≠ 0
    using norm-and-cross-eq-0[of n y - x] assms(3) xy-neq
    by simp
  have ?a · y = ?b
    by (metis (no-types, opaque-lifting) add.commute diff-0-right diff-add-cancel
      dot-cross-self(2) inner-commute
      inner-diff-left)

  have F exposed-face-of (convex hull S)
    using *(1) polyhedron-convex-hull[OF assms(2)] exposed-face-of-polyhedron
    by blast
  then obtain a0 b0 where a0-b0-facts:
    convex hull S ⊆ {w. a0 · w ≤ b0}
    convex hull T = convex hull S ∩ {w. a0 · w = b0}
    unfolding exposed-face-of-def T-facts(2) by blast
  define a' where a' = a0 - ((a0 · n) / (n · n)) *R n
  define b' where b' = b0 - ((a0 · n) / (n · n)) * d
  have a' · n = 0

```

```

    unfolding a'-def
    by (simp add: inner-diff-left dot-square-norm power2-eq-square)
have  $S \subseteq \{w. n \cdot w = d\}$ 
    using assms(1) by blast
have  $\bigwedge w. w \in \text{convex hull } S \implies (a' \cdot w = b') = (a0 \cdot w = b0)$ 
proof -
    fix w assume  $w \in \text{convex hull } S$ 
    then have  $n \cdot w = d$ 
        using  $\langle S \subseteq \{w. n \cdot w = d\} \rangle$ 
        by (metis (mono-tags, lifting) assms(1) convex-hyperplane hull-induct)
    then show  $(a' \cdot w = b') = (a0 \cdot w = b0)$ 
        unfolding a'-def b'-def by (simp add: inner-diff-left)
qed
moreover have  $\bigwedge w. w \in \text{convex hull } S \implies (a' \cdot w \leq b') = (a0 \cdot w \leq b0)$ 
proof -
    fix w assume  $w \in \text{convex hull } S$ 
    then have  $n \cdot w = d$ 
        using  $\langle S \subseteq \{w. n \cdot w = d\} \rangle$ 
        by (metis (mono-tags, lifting) assms(1) convex-hyperplane hull-induct)
    then show  $(a' \cdot w \leq b') = (a0 \cdot w \leq b0)$ 
        unfolding a'-def b'-def by (simp add: inner-diff-left)
qed
ultimately have a'-b'-facts:
    convex hull  $S \subseteq \{w. a' \cdot w \leq b'\}$ 
    convex hull  $T = \text{convex hull } S \cap \{w. a' \cdot w = b'\}$ 
    using a0-b0-facts by auto
let ?H' =  $\{w. a' \cdot w = b'\}$ 

have aff-dim (affine hull  $T$ ) = 1
    by (metis U-facts(3)  $\langle \text{aff-dim } U = 1 \rangle$  aff-dim-affine-hull)
have  $\{x, y\} \subseteq ?H$ 
    using  $\langle ?a \cdot y = ?b \rangle$  by simp
then have affine hull  $\{x, y\} \subseteq ?H$ 
    using affine-hyperplane hull-minimal by metis
then have affine hull  $T \subseteq ?H$ 
    using U-eq U-facts(3) by blast
moreover have affine hull  $T \subseteq ?H'$ 
    using a'-b'-facts(2)
    by (metis affine-hyperplane affine-imp-convex inf.cobounded2 subset-hull)
ultimately have affine hull  $T \subseteq ?H \cap ?H'$ 
    by blast
then have aff-dim  $(?H \cap ?H') \geq 1$ 
    using aff-dim-subset  $\langle \text{aff-dim } (\text{affine hull } T) = 1 \rangle$ 
    by metis

show ?rhs
proof (cases convex hull  $T = \text{convex hull } S$ )
case True
    then have affine hull  $T = \text{affine hull } S$ 

```

```

    using affine-hull-convex-hull by metis
  then have  $S \subseteq \text{affine hull } T$ 
    using T-facts(1)
    by (simp add: hull-subset)
  then have  $S \subseteq ?H$ 
    using  $\langle \text{affine hull } T \subseteq ?H \rangle$  by blast
  then show ?thesis
    using xy-in-S  $\langle ?a \neq 0 \rangle$  True T-facts(2) unfolding Let-def
    by (metis (mono-tags, lifting) Int-Collect dual-order.refl inf.orderE)

```

next

```

  case False
  have affine hull  $T \subseteq ?H'$ 
    using a'-b'-facts(2) T-facts(2)
    by (metis affine-hyperplane affine-imp-convex inf.cobounded2 subset-hull)

```

```

  have affine hull  $T \subseteq ?H \cap ?H'$ 
    using  $\langle \text{affine hull } T \subseteq ?H \rangle \langle \text{affine hull } T \subseteq ?H' \rangle$  by blast

```

```

  have aff-dim  $?H = 2$ 
    using  $\langle ?a \neq 0 \rangle$  aff-dim-hyperplane by simp
  have  $a' \neq 0$ 
    using *(2) False Int-commute T-facts(2) a'-b'-facts(2)
    by auto
  then have aff-dim  $?H' = 2$ 
    using aff-dim-hyperplane by simp

```

```

  have  $?H \subseteq ?H'$ 
  proof (rule ccontr)
    assume  $\neg(?H \subseteq ?H')$ 
    then have aff-dim  $(?H \cap ?H') = 1$ 
      using aff-dim-affine-Int-hyperplane[of ?H a' b']
       $\langle \text{aff-dim } ?H = 2 \rangle \langle \text{aff-dim } ?H' = 2 \rangle \langle \text{aff-dim } (?H \cap ?H') \geq 1 \rangle$ 
    by (metis (no-types, lifting) ext affine-hyperplane diff-add-cancel diff-minus-eq-add
      diff-numeral-special(11) not-numeral-le-neg-one numeral-eq-one-iff)
    have  $x + n \in ?H$ 
      by (simp add: dot-cross-self(4) inner-right-distrib)
    have  $x \in ?H'$ 
      using U-eq U-facts(1) a'-b'-facts(2) hull-inc by fastforce
    have  $n \cdot (x + n) \neq d$ 
      by (simp add: assms(1,3) inner-right-distrib xy-in-S(1))
    have affine hull  $S \subseteq \{w. n \cdot w = d\}$ 
      by (simp add:  $\langle S \subseteq \{w. n \cdot w = d\} \rangle$  affine-hyperplane subset-hull)
    then have convex hull  $T \subseteq \{w. n \cdot w = d\}$ 
      using a'-b'-facts(2) convex-hull-subset-affine-hull by auto
    have  $x + n \in ?H'$ 
      using  $\langle a' \cdot n = 0 \rangle$ 
      by (metis  $\langle x \in ?H' \rangle$  add commute add-0 inner-right-distrib mem-Collect-eq)
    have  $?H \cap ?H' \subseteq \{w. n \cdot w = d\}$ 

```

```

    by (metis (mono-tags, lifting) ⟨aff-dim (affine hull T) = 1⟩ ⟨aff-dim (?H ∩
    ?H') = 1⟩
    ⟨affine hull T ⊆ ?H ∩ ?H'⟩ ⟨convex hull T ⊆ {w. n • w = d}⟩
    aff-dim-eq-full-gen
    affine-hull-convex-hull affine-hyperplane subset-hull)
  then show False
    using ⟨x + n ∈ ?H⟩ ⟨x + n ∈ ?H'⟩ ⟨n • (x + n) ≠ d⟩
    by blast
qed
then have ?H = ?H'
  using subset-hyperplanes
  by (metis ⟨a' ≠ 0⟩ ⟨{x, y} ⊆ ?H⟩ empty-iff hyperplane-eq-UNIV insert-subset)

then have F = convex hull (S ∩ ?H)
  by (smt (verit, ccfv-threshold) Int-commute T-facts(1,2) a'-b'-facts(2) con-
  vex-convex-hull
  hull-minimal hull-subset inf.orderE le-inf-iff subsetI)

moreover have (∀ w ∈ S. ?a • w ≤ ?b) ∨ (∀ w ∈ S. ?a • w ≥ ?b)
proof -
  have ?H' ⊆ ?H
    using ⟨?H = ?H'⟩ by auto
  then obtain k where k-facts: ?a = k *R a' ?b = k * b' k ≠ 0
    using ⟨?a ≠ 0⟩ ⟨a' ≠ 0⟩ hyperplane-subset-imp-scaled by blast
  show ?thesis
proof (cases k > 0)
  case True
    then have ∧w. ?a • w ≤ ?b ⟷ a' • w ≤ b'
      using k-facts by (auto simp: mult-le-cancel-left)
    then show ?thesis
      using a'-b'-facts(1) hull-subset[of S] by auto
  next
  case False
    then have k < 0
      using k-facts by simp
    then have ∧w. ?a • w ≥ ?b ⟷ a' • w ≤ b'
      using k-facts by (auto simp: mult-le-cancel-left)
    then show ?thesis
      using a'-b'-facts(1) hull-subset[of S] by auto
qed
qed

ultimately show ?thesis
  using xy-in-S ⟨?a ≠ 0⟩ unfolding Let-def
  by blast
qed
next
assume ?rhs
then obtain x y where

```

```

xy-in-S:  $x \in S \ y \in S$  and
a-neq-0:  $n \times (y - x) \neq 0$  and
one-side:  $(\forall w \in S. n \times (y - x) \cdot w \leq n \times (y - x) \cdot x) \vee$ 
            $(\forall w \in S. n \times (y - x) \cdot w \geq n \times (y - x) \cdot x)$  and
F-eq:  $F = \text{convex hull } (S \cap \{w. n \times (y - x) \cdot w = n \times (y - x) \cdot x\})$ 
unfolding Let-def by blast
let ?a =  $n \times (y - x)$ 
let ?b = ?a · x
let ?H =  $\{w. ?a \cdot w = ?b\}$ 
have convex hull  $(S \cap ?H) = \text{convex hull } S \cap ?H$ 
proof
  show convex hull  $(S \cap ?H) \subseteq \text{convex hull } S \cap ?H$ 
  using convex-hyperplane[of ?a ?b] hull-mono[of  $S \cap ?H$  S convex]
  subset-hull[of convex ?H S  $\cap ?H$ ]
  by blast
next
show convex hull  $S \cap ?H \subseteq \text{convex hull } (S \cap ?H)$ 
proof (cases  $S \subseteq ?H$ )
  case True
  then show ?thesis by (simp add: Int-absorb2)
next
  case False
  let ?S-on =  $S \cap ?H$ 
  let ?S-off =  $S - ?H$ 
  have convex hull  $S \cap ?H = \text{convex hull } (\text{convex hull } (?S\text{-on}) \cup \text{convex hull } (?S\text{-off})) \cap ?H$ 
  by (metis Int-Diff-Un hull-Un-left hull-Un-right)
  moreover have convex hull  $(\text{convex hull } (?S\text{-on}) \cup \text{convex hull } (?S\text{-off})) \cap ?H \subseteq \text{convex hull } (S \cap ?H)$ 
  proof
    fix w
    assume *:  $w \in \text{convex hull } (\text{convex hull } (?S\text{-on}) \cup \text{convex hull } (?S\text{-off})) \cap ?H$ 
    have convex: convex  $(\text{convex hull } (?S\text{-on}))$  convex  $(\text{convex hull } (?S\text{-off}))$ 
    using convex-convex-hull by blast+
    have nonempty1:  $\text{convex hull } ?S\text{-on} \neq \{\}$ 
    using xy-in-S(1) convex-hull-eq-empty by blast
    have nonempty2:  $\text{convex hull } ?S\text{-off} \neq \{\}$ 
    using False convex-hull-eq-empty by simp
    obtain u p q where
      w-facts:  $w = (1 - u) *_R p + u *_R q$  ?a · w = ?b and
      p-in:  $p \in \text{convex hull } ?S\text{-on}$  and
      q-in:  $q \in \text{convex hull } ?S\text{-off}$  and
      u-range:  $0 \leq u \leq 1$ 
    using * convex-hull-union-nonempty-explicit[OF convex(1) nonempty1 convex(2) nonempty2]
    by blast
    show  $w \in \text{convex hull } (S \cap ?H)$ 
    proof (cases  $u = 0$ )

```

```

case True
then show ?thesis
  using p-in w-facts(1) by simp
next
case u-nonzero: False
have  $p \in \text{convex hull } ?H$ 
  using p-in hull-mono[of ?S-on ?H] by blast
then have  $?a \cdot p = ?b$ 
  using convex-hyperplane hull-same
  by blast
have  $q \in \text{convex hull } ?S\text{-off}$ 
  using q-in by simp

have  $q \notin ?H$ 
proof
  assume  $q \in ?H$ 
  then have  $?a \cdot q = ?b$  by simp
  consider
    (leq)  $\forall w \in S. ?a \cdot w \leq ?b$  |
    (geq)  $\forall w \in S. ?a \cdot w \geq ?b$ 
  using one-side by blast
  then show False
  proof cases
    case leq
    then have  $?S\text{-off} \subseteq \{w. ?a \cdot w < ?b\}$ 
    by fastforce
    then have  $\text{convex hull } ?S\text{-off} \subseteq \{w. ?a \cdot w < ?b\}$ 
    using convex-halfspace-lt hull-minimal by metis
    then show False
    using  $\langle ?a \cdot q = ?b \rangle \langle q \in \text{convex hull } ?S\text{-off} \rangle$  by auto
  next
    case geq
    then have  $?S\text{-off} \subseteq \{w. ?a \cdot w > ?b\}$ 
    by fastforce
    then have  $\text{convex hull } ?S\text{-off} \subseteq \{w. ?a \cdot w > ?b\}$ 
    using convex-halfspace-gt hull-minimal by metis
    then show False
    using  $\langle ?a \cdot q = ?b \rangle \langle q \in \text{convex hull } ?S\text{-off} \rangle$  by auto
  qed
qed
then have q-not-in-H: ?a · q ≠ ?b by simp

have  $?a \cdot ((1 - u) *_R p + u *_R q) = (1 - u) * (?a \cdot p) + u * (?a \cdot q)$ 
  by (simp add: inner-add-right)
then show ?thesis
  using q-not-in-H  $\langle q \in \text{convex hull } ?S\text{-off} \rangle$ 
  by (metis  $\langle ?a \cdot p = ?b \rangle$  real-scaleR-def segment-degen-0 u-nonzero w-facts)
qed
qed

```

ultimately show *?thesis* by *argo*
 qed
 qed
 then have *F-as-inter*: $F = \text{convex hull } S \cap ?H$
 using *F-eq* by *simp*
 consider (*case-le*) $\forall w \in S. ?a \cdot w \leq ?b$ | (*case-ge*) $\forall w \in S. ?a \cdot w \geq ?b$
 using *one-side* by *blast*
 then have *face-of*: $F \text{ face-of } (\text{convex hull } S)$
 proof *cases*
 case *case-le*
 then have *subset*: $S \subseteq \{x. ?a \cdot x \leq ?b\}$
 by *blast*
 have $\text{convex hull } S \subseteq \{x. ?a \cdot x \leq ?b\}$
 using *convex-halfspace-le*[*of ?a ?b*] *hull-minimal*[*OF subset, of convex*]
 by *blast*
 then have $\bigwedge w. w \in \text{convex hull } S \implies ?a \cdot w \leq ?b$
 by *blast*
 then have $\text{convex hull } S \cap ?H \text{ face-of } \text{convex hull } S$
 using *face-of-Int-supporting-hyperplane-le*[*OF convex-convex-hull*[*of S*]]
 by *blast*
 then show *?thesis*
 using *F-as-inter* by *simp*
 next
 case *case-ge*
 then have *subset*: $S \subseteq \{x. ?a \cdot x \geq ?b\}$
 by *blast*
 have $\text{convex hull } S \subseteq \{x. ?a \cdot x \geq ?b\}$
 using *convex-halfspace-ge*[*of ?b ?a*] *hull-minimal*[*OF subset, of convex*]
 by *blast*
 then have $\bigwedge w. w \in \text{convex hull } S \implies ?a \cdot w \geq ?b$
 by *blast*
 then have $\text{convex hull } S \cap ?H \text{ face-of } \text{convex hull } S$
 using *face-of-Int-supporting-hyperplane-ge*[*OF convex-convex-hull*[*of S*]]
 by *blast*
 then show *?thesis*
 using *F-as-inter* by *simp*
 qed
 have $S \subseteq \{x. n \cdot x = d\}$
 using *assms*(1) by *blast*
 have *aff-dim* $?H = 2$
 using *aff-dim-hyperplane*[*OF a-neq-0, of ?b*] by *simp*
 have $\neg(?H \subseteq \{v. n \cdot v = d\})$
 proof –
 have $\exists v. n \cdot v \neq d \wedge n \times (y - x) \cdot v = x \cdot n \times (y - x)$
 by (*metis* (*no-types*) *assms*(3) *diff-0-right* *diff-add-cancel* *dot-cross-self*(1))
inner-commute
 inner-diff-left inner-eq-zero-iff minus-diff-eq)
 then show *?thesis*

```

    by (smt (z3) Collect-mono-iff inner-commute)
qed
then have aff-dim (?H ∩ {x. n • x = d}) ≤ 1
  using aff-dim-affine-Int-hyperplane[of ?H n d] unfolding ⟨aff-dim ?H = 2⟩
  by (simp add: affine-hyperplane)
moreover have F ⊆ (?H ∩ {x. n • x = d})
  using F-as-inter assms(1) ⟨S ⊆ {x. n • x = d}⟩
    convex-hyperplane[of n d] subset-hull[of convex {x. n • x = d} S]
  by blast
ultimately have upper: aff-dim F ≤ 1
  using aff-dim-subset[of F ?H ∩ {x. n • x = d}] unfolding ⟨aff-dim ?H = 2⟩
  by linarith
have x ≠ y
  using a-neq-0 by auto
moreover have ?a • y = ?b
  by (metis (no-types, opaque-lifting) dot-cross-self(2) inner-commute inner-diff-left
right-minus-eq)
moreover have {x, y} ⊆ F
  using xy-in-S
  by (simp add: F-as-inter calculation(2) hull-inc)
moreover have aff-dim {x, y} = 1
  by (simp add: calculation(1))
ultimately have lower: aff-dim F ≥ 1
  using aff-dim-subset by metis

show ?lhs
  using face-of upper lower by simp
qed

lemma compute-faces-1-step:
  fixes F n u S and T :: (real^3) set
  defines Q ≡ (λx y. let a = n × (y - x) in
    a ≠ 0 ∧
    (let b = a • x in
      ((∀ w ∈ T. a • w ≤ b) ∨ (∀ w ∈ T. a • w ≥ b)) ∧
      F = convex hull (T ∩ {k. a • k = b})))
  shows (∃ y ∈ (insert u S). ∃ z ∈ (insert u S). Q y z) =
    ((∃ z ∈ S. Q u z) ∨ (∃ y ∈ S. ∃ z ∈ S. Q y z))
  (is ?lhs = ?rhs)
proof -
  have Q-degen: ∧x. ¬ Q x x
    unfolding Q-def Let-def
    by (simp add: cross3-simps)

  have Q-sym: ∧x y. Q x y ⟹ Q y x
  proof -
    fix x y
    assume Q x y
    then obtain a b where

```

```

    a-def:  $a = n \times (y - x)$  and
    a-nonzero:  $a \neq 0$  and
    b-def:  $b = a \cdot x$  and
    ineq:  $(\forall w \in T. a \cdot w \leq b) \vee (\forall w \in T. a \cdot w \geq b)$  and
    plane:  $F = \text{convex hull } (T \cap \{w. a \cdot w = b\})$ 
    unfolding Q-def Let-def
    by auto
  let ?a' =  $n \times (x - y)$ 
  let ?b' =  $?a' \cdot y$ 
  have ?a' =  $-a$ 
    unfolding a-def Cross3.right-diff-distrib
    by auto
  have  $a \cdot y - a \cdot x = 0$ 
    by (metis a-def dot-cross-self(3) inner-diff-right)
  then have  $\{w. ?a' \cdot w = ?b'\} = \{w. a \cdot w = b\}$ 
    unfolding ⟨?a' = -a⟩ b-def
    by auto
  then show Q y x
    unfolding Q-def Let-def
    using ⟨?a' = -a⟩ a-nonzero plane ineq
    by auto
qed

show ?thesis
  using set-choice-decomposition-2[of Q u S] Q-sym Q-degen Q-def
  by blast
qed

method facet-edges uses plane n-nz defs =
  (simp only: compute-faces-1[OF plane - n-nz] finite-insert finite.emptyI,
   simp only: compute-faces-1-step bex-empty simp-thms(31),
   simp only: empty-iff ball-insert bex-empty bex-simps(5) ball-simps(5) simp-thms(21,31),
   simp add: defs vector3-sub vector3-cross vector3-dot vector3-eq-0 Let-def
            rat5-mul rat5-add rat5-sub rat5-eq rat5-eq-0 rat5-le)

```

4.4 Full-dimensionality

```

lemma polytope-3D:
  fixes S :: (real^3) set
  assumes  $(z - x) \times (y - x) \neq 0$ 
    and  $\exists w \in S. ((z - x) \times (y - x)) \cdot w \neq ((z - x) \times (y - x)) \cdot x$ 
  shows  $\text{aff-dim } (\text{convex hull } (\text{insert } x (\text{insert } y (\text{insert } z S)))) = 3$ 
    (is  $\text{aff-dim } ?\text{hull} = 3$ )
proof -
  have  $\text{aff-dim } ?\text{hull} \leq \text{int } (\text{DIM } (\text{real}^3))$ 
    by (rule aff-dim-le-DIM)
  then have upper:  $\text{aff-dim } ?\text{hull} \leq 3$ 
    by auto
  let ?a =  $(z - x) \times (y - x)$ 

```

```

obtain  $w$  where  $w$ -facts:  $w \in S \text{ ?}a \cdot w \neq \text{?}a \cdot x$ 
  using  $assms(2)$  by  $blast$ 
have  $\{w, x, y, z\} \subseteq \text{?}hull$ 
  using  $hull$ -subset  $w$ -facts(1) by  $fastforce$ 
have  $aff$ -dim  $\{w, x, y, z\} \geq 3$ 
proof (cases  $w \in affine\ hull\ \{x, y, z\}$ )
  case  $True$ 
    have  $\text{?}a \cdot y = \text{?}a \cdot x \text{ ?}a \cdot z = \text{?}a \cdot x$ 
    by ( $metis$  ( $no$ -types,  $lifting$ )  $Cross3$ .left-diff-distrib  $cross$ -refl  $cross$ -skew  $cross$ -triple
       $diff$ -0  $diff$ -0-right)+
    then have  $\{x, y, z\} \subseteq \{v. \text{?}a \cdot v = \text{?}a \cdot x\}$ 
    by  $blast$ 
    then have  $w \in \{v. \text{?}a \cdot v = \text{?}a \cdot x\}$ 
    using  $True$   $affine$ -hyperplane[of  $\text{?}a \text{ ?}a \cdot x$ ]
       $subset$ -hull[of  $affine\ \{v. \text{?}a \cdot v = \text{?}a \cdot x\}\ \{x, y, z\}$ ]
    by  $blast$ 
    then show  $\text{?thesis}$ 
    using  $w$ -facts(2) by  $blast$ 
  next
    case  $False$ 
    have  $\neg collinear\ \{x, y, z\}$ 
    by ( $metis$  ( $no$ -types,  $opaque$ -lifting)  $assms(1)$   $NO$ -MATCH-def  $collinear$ -3
       $cross$ -eq-0
       $insert$ -commute)
    then show  $\text{?thesis}$ 
    using  $False$   $aff$ -dim-insert
    by ( $smt$  ( $verit$ ,  $ccfv$ -threshold)  $collinear$ - $aff$ -dim)
  qed
then have  $aff$ -dim  $\text{?}hull \geq 3$ 
  by ( $meson$  ( $\{w, x, y, z\} \subseteq \text{?}hull$ )  $aff$ -dim-subset  $order$ -trans)
then show  $\text{?thesis}$ 
  using  $upper$  by  $simp$ 
qed

```

4.5 Relations between faces of different dimensions

These lemmas are used when formalizing the number of edges per face and number of edges per vertex.

corollary $edge$ -belongs-to-face:

```

assumes  $e$  face-of  $p$ 
  and  $aff$ -dim  $e = 1$ 
  and  $aff$ -dim  $p = 3$ 
  and  $polyhedron\ p$ 
obtains  $f$  where  $f$  face-of  $p$   $aff$ -dim  $f = 2$   $e$  face-of  $f$ 
using  $assms(2,3)$   $ridge$ -belongs-to-facet[ $OF\ assms(1) - - assms(4)$ ]
by  $auto$ 

```

corollary $vertex$ -belongs-to-edge:

```

assumes  $v$  face-of  $p$ 

```

and *aff-dim* $v = 0$
and *aff-dim* $p = 3$
and *polyhedron* p
obtains e **where** e *face-of* p *aff-dim* $e = 1$ v *face-of* e
using *assms*(2,3) *peak-belongs-to-ridge*[*OF assms*(1) - - *assms*(4)]
by *auto*

lemma *vertices-of-edge*:
assumes $e = \text{convex hull } \{a, b\}$
and v *face-of* e
and *aff-dim* $v = 0$
shows $v = \{a\} \vee v = \{b\}$
by (*metis assms aff-dim-eq-0 extreme-point-of-convex-hull-2 face-of-singleton*)

lemma *vertex-edge-set-eq*:
assumes v *face-of* S
shows $\text{card } \{e. e \text{ face-of } S \wedge \text{aff-dim } e = 1 \wedge v \subseteq e\} =$
 $\text{card } \{e. e \text{ face-of } S \wedge \text{aff-dim } e = 1 \wedge v \text{ face-of } e\}$
by (*meson assms face-of-face*)

lemma *face-edge-set-eq*:
assumes f *face-of* S
shows $\text{card } \{e. e \text{ face-of } S \wedge \text{aff-dim } e = 1 \wedge e \subseteq f\} =$
 $\text{card } \{e. e \text{ face-of } f \wedge \text{aff-dim } e = 1\}$
using *face-of-face*[*OF assms*(1)]
by *meson*

lemma *set-disj-union-distr*:
shows $\{e. (P e \vee Q e) \wedge R e\} = \{e. P e \wedge R e\} \cup \{e. Q e \wedge R e\}$
by *blast*

lemma *v-sing-face-of-e*:
shows $(e = a \wedge \{x\} \text{ face-of } e) = (e = a \wedge x \text{ extreme-point-of } a)$
using *face-of-singleton* **by** *blast*

lemma *card-3-iff2*:
shows $(\text{card } \{x, y, z\} = 3) = (x \neq y \wedge x \neq z \wedge y \neq z)$
by (*simp add: card-insert-if*)

lemma *card-4-iff2*:
shows $(\text{card } \{w, x, y, z\} = 4) = (w \neq x \wedge w \neq y \wedge w \neq z \wedge x \neq y \wedge x \neq z \wedge$
 $y \neq z)$
by (*simp add: card-insert-if*)

lemma *card-5-iff*:
shows $\text{card } \{u, w, x, y, z\} = 5 \longleftrightarrow (u \neq w \wedge u \neq x \wedge u \neq y \wedge u \neq z \wedge w \neq x \wedge$
 $w \neq y \wedge w \neq z \wedge x \neq y \wedge y \neq z \wedge x \neq z)$
by (*simp add: card-insert-if*)

method *edges-per-vertex* =
 (unfold *set-disj-union-distr* *segment-convex-hull*[*symmetric*] *v-sing-face-of-e*;
simp-all add: *extreme-point-of-segment* *vector3-eq* *rat5-eq*;
 unfold *card-3-iff2* *card-4-iff2* *card-5-iff* *closed-segment-eq*;
simp-all add: *vector3-eq* *rat5-eq* *doubleton-eq-iff*)

lemma *set-cases*:
shows $\{x. x = a\} = \{a\} \ \{x. x = a \vee P \ x\} = \text{insert } a \ \{x. P \ x\}$
by *auto*

4.6 Congruence

definition *congruent* :: [*'a::euclidean-space set*, *'a set*] $\Rightarrow$ bool (**infixr** $\langle$ (*congruent*) $\rangle$
 50)

where *S congruent T* $\longleftrightarrow (\exists c \ f. \text{orthogonal-transformation } f \wedge T = (\lambda x. c + f \ x) \ 'S)$

lemma *congruent-refl*:
shows *S congruent S*
unfolding *congruent-def*
using *orthogonal-transformation-id* *image-add-0*
by *blast*

lemma *congruent-sym*:
shows *S congruent T* $\longleftrightarrow T \text{ congruent } S$
proof –
have *forward*: $\bigwedge S \ T. S \text{ congruent } T \implies T \text{ congruent } S$
proof –
fix *S T* :: *'b set*
assume *S congruent T*
then obtain *c f* **where** *orthog*: *orthogonal-transformation f* **and** *T*: $T = (\lambda x. c + f \ x) \ 'S$
unfolding *congruent-def* **by** *blast*
then have *orthog-inv*: *orthogonal-transformation (inv f)*
using *orthogonal-transformation-inv* **by** *blast*
have *bounded-linear (inv f)*
using *orthogonal-transformation-linear*[*OF orthog-inv*]
by (*simp* add: *linear-conv-bounded-linear*)

have $(\lambda y. - \text{inv } f \ c + \text{inv } f \ y) \ 'T = (\lambda x. - \text{inv } f \ c + \text{inv } f \ (c + f \ x)) \ 'S$
using *T* **by** *blast*
also have $\dots = (\lambda x. - \text{inv } f \ c + \text{inv } f \ c + \text{inv } f \ (f \ x)) \ 'S$
using $\langle$ *bounded-linear (inv f)* $\rangle$
by (*simp* add: *linear-simps*)
also have $\dots = (\lambda x. \text{inv } f \ (f \ x)) \ 'S$
by *simp*
finally have $S = (\lambda y. - \text{inv } f \ c + \text{inv } f \ y) \ 'T$
by (*simp* add: *orthog bijection.intro bijection.inv-left orthogonal-transformation-bij*)
then show *T congruent S*

```

    unfolding congruent-def using orthog-inv
  by blast
qed
then show ?thesis
  by blast
qed

lemma congruent-trans:
  assumes  $R$  congruent  $S$ 
  assumes  $S$  congruent  $T$ 
  shows  $R$  congruent  $T$ 
proof -
  obtain  $f$   $c$  where  $orthog-f$ : orthogonal-transformation  $f$  and  $S$ :  $S = (\lambda x. c + f$ 
 $x)$  ‘  $R$ 
    using assms(1) unfolding congruent-def by blast
  obtain  $g$   $d$  where  $orthog-g$ : orthogonal-transformation  $g$  and  $T$ :  $T = (\lambda x. d +$ 
 $g$   $x)$  ‘  $S$ 
    using assms(2) unfolding congruent-def by blast
  then have  $T = (\lambda x. d + g (c + f x))$  ‘  $R$ 
    unfolding  $S$  by blast
  then have  $T = (\lambda x. d + g c + g (f x))$  ‘  $R$ 
    using orthogonal-transformation-linear[ $OF$   $orthog-g$ ] linear-conv-bounded-linear[ $of$ 
 $g$ ]
    linear-simps(1)
  by fastforce
  moreover have orthogonal-transformation  $(g \circ f)$ 
    using  $orthog-f$   $orthog-g$  orthogonal-transformation-compose
  by blast
  ultimately show ?thesis
    unfolding congruent-def
  by fastforce
qed

lemma congruent-set:
  assumes  $\forall t \in (insert\ s\ S). s$  congruent  $t$ 
  assumes  $x \in (insert\ s\ S)$ 
  assumes  $y \in (insert\ s\ S)$ 
  shows  $x$  congruent  $y$ 
  using assms congruent-sym congruent-trans
  by metis

```

4.6.1 Congruence of segments

```

lemma congruent-segments:
  fixes  $a$   $b$   $c$   $d :: real^n$ 
  assumes  $dist\ a\ b = dist\ c\ d$ 
  shows  $(closed-segment\ a\ b)$  congruent  $(closed-segment\ c\ d)$ 
proof -
  have  $norm\ (b - a) = norm\ (d - c)$ 

```

```

using assms
by (simp add: dist-norm norm-minus-commute)
then obtain f where f-facts: orthogonal-transformation f f (b - a) = d - c
using orthogonal-transformation-exists
by blast
have linear f
using orthogonal-transformation-linear[OF f-facts(1)] .
have  $d = (c - f\ a) + f\ b$ 
using linear-diff[OF ⟨linear f⟩, of b a] unfolding f-facts(2)
by (metis add.left-commute diff-add-cancel)
moreover have  $(\lambda x. (c - f\ a) + f\ x) = (\lambda x. (c - f\ a) + x) \circ f$ 
by auto
ultimately have  $(\lambda x. (c - f\ a) + f\ x) \text{ ‘ closed-segment } a\ b =$ 
 $((\lambda x. (c - f\ a) + x) \circ f) \text{ ‘ (convex hull } \{a, b\})$ 
unfolding segment-convex-hull[of a b]
by arg0
also have  $\dots = (\lambda x. (c - f\ a) + x) \text{ ‘ (convex hull (f ‘ \{a, b\}))}$ 
using image-comp convex-hull-linear-image[OF ⟨linear f⟩]
by metis
also have  $\dots = \text{convex hull } \{(c - f\ a) + f\ a, (c - f\ a) + f\ b\}$ 
by (metis (no-types, opaque-lifting) convex-hull-translation
image-empty image-insert)
also have  $\dots = \text{convex hull } \{c, d\}$ 
unfolding  $\langle d = (c - f\ a) + f\ b \rangle$  by auto
finally show ?thesis
unfolding congruent-def segment-convex-hull[of c d]
using f-facts(1)
by blast
qed

```

4.6.2 Congruence of 2-faces

lemma *congruent-simple*:

assumes $\exists A::\text{real}^{\mathcal{I}\mathcal{I}}. \text{orthogonal-matrix } A \wedge (\lambda x. A * v\ x) \text{ ‘ } S = T$
shows $(\text{convex hull } S) \text{ congruent } (\text{convex hull } T)$

proof –

```

obtain A where A-facts: orthogonal-matrix A (λx. A * v x) ‘ S = T
using assms by blast
then have convex hull T = (λx. A * v x) ‘ (convex hull S)
using convex-hull-linear-image by blast
moreover have orthogonal-transformation (λx. A * v x)
using A-facts(1) orthogonal-transformation-matrix
by fastforce
ultimately show ?thesis
unfolding congruent-def
by (metis image-add-0[of convex hull T]
image-image[of (+) 0 (*v) A convex hull S])

```

qed

lemma *matrix-vector-mul-3*:

shows (*vector* [*vector* [*a11*, *a12*, *a13*],
vector [*a21*, *a22*, *a23*],
vector [*a31*, *a32*, *a33*]]::*real*³³) **v* *vector* [*x1*, *x2*, *x3*] =
vector [*a11* * *x1* + *a12* * *x2* + *a13* * *x3*,
a21 * *x1* + *a22* * *x2* + *a23* * *x3*,
a31 * *x1* + *a32* * *x2* + *a33* * *x3*]
by (*simp add: vec-eq-iff matrix-vector-mult-def forall-3 sum-3*)

lemma *matrix-lemma*:

fixes *A* :: *real*³³
shows (*A* **v* *x1* = *x2* ∧ *A* **v* *y1* = *y2* ∧ *A* **v* *z1* = *z2*) =
((*vector* [*x1*, *y1*, *z1*]::*real*³³) **v* (*row* 1 *A*) = *vector* [*x2*\$1, *y2*\$1, *z2*\$1] ∧
(*vector* [*x1*, *y1*, *z1*]::*real*³³) **v* (*row* 2 *A*) = *vector* [*x2*\$2, *y2*\$2, *z2*\$2] ∧
(*vector* [*x1*, *y1*, *z1*]::*real*³³) **v* (*row* 3 *A*) = *vector* [*x2*\$3, *y2*\$3, *z2*\$3])
by (*simp add: vec-eq-iff matrix-vector-mult-def row-def forall-3 sum-3 mult.commute conj-comms*)

lemma *matrix-by-cramer-lemma*:

fixes *A* :: *real*³³
assumes *det* (*vector* [*x1*, *y1*, *z1*]::*real*³³) ≠ 0
shows (*A* **v* *x1* = *x2* ∧ *A* **v* *y1* = *y2* ∧ *A* **v* *z1* = *z2*) =
(*A* = (χ *m k*. *det* ((χ *i j*.
if *j* = *k*
then (*vector* [*x2*\$*m*, *y2*\$*m*, *z2*\$*m*]::*real*³³)\$*i*
else (*vector* [*x1*, *y1*, *z1*]::*real*³³)\$*i*\$*j*)
::*real*³³) /
det (*vector* [*x1*, *y1*, *z1*]::*real*³³)))
unfolding *matrix-lemma cramer*[*OF* *assms*(1)]
by (*simp add: vec-eq-iff row-def forall-3*)

definition *mk-matrix33* :: (*real*³) ⇒ (*real*³) ⇒ (*real*³) ⇒ (*real*³) ⇒ (*real*³)
⇒ (*real*³) ⇒ *real*³³

where *mk-matrix33* *x1 y1 z1 x2 y2 z2* = (*let* *d* = *det* (*vector* [*x1*, *y1*, *z1*]::*real*³³)
in *vector* [*vector* [*vector* [*vector* [*x2*\$1 * *y1*\$2 * *z1*\$3 +

x1\$2 * *y1*\$3 * *z2*\$1 +
x1\$3 * *y2*\$1 * *z1*\$2 −
x2\$1 * *y1*\$3 * *z1*\$2 −
x1\$2 * *y2*\$1 * *z1*\$3 −
x1\$3 * *y1*\$2 * *z2*\$1) / *d*,
(*vector* [*vector* [*vector* [*x1*\$1 * *y2*\$1 * *z1*\$3 +
x2\$1 * *y1*\$3 * *z1*\$1 +
x1\$3 * *y1*\$1 * *z2*\$1 −
x1\$1 * *y1*\$3 * *z2*\$1 −
x2\$1 * *y1*\$1 * *z1*\$3 −
x1\$3 * *y2*\$1 * *z1*\$1) / *d*,
(*vector* [*vector* [*vector* [*x1*\$1 * *y1*\$2 * *z2*\$1 +
x1\$2 * *y2*\$1 * *z1*\$1 +

$$\begin{aligned}
& x2\$1 * y1\$1 * z1\$2 - \\
& x1\$1 * y2\$1 * z1\$2 - \\
& x1\$2 * y1\$1 * z2\$1 - \\
& x2\$1 * y1\$2 * z1\$1) / d], \text{ vector } [\\
& (x2\$2 * y1\$2 * z1\$3 + \\
& x1\$2 * y1\$3 * z2\$2 + \\
& x1\$3 * y2\$2 * z1\$2 - \\
& x2\$2 * y1\$3 * z1\$2 - \\
& x1\$2 * y2\$2 * z1\$3 - \\
& x1\$3 * y1\$2 * z2\$2) / d, \\
& (x1\$1 * y2\$2 * z1\$3 + \\
& x2\$2 * y1\$3 * z1\$1 + \\
& x1\$3 * y1\$1 * z2\$2 - \\
& x1\$1 * y1\$3 * z2\$2 - \\
& x2\$2 * y1\$1 * z1\$3 - \\
& x1\$3 * y2\$2 * z1\$1) / d, \\
& (x1\$1 * y1\$2 * z2\$2 + \\
& x1\$2 * y2\$2 * z1\$1 + \\
& x2\$2 * y1\$1 * z1\$2 - \\
& x1\$1 * y2\$2 * z1\$2 - \\
& x1\$2 * y1\$1 * z2\$2 - \\
& x2\$2 * y1\$2 * z1\$1) / d], \text{ vector } [\\
& (x2\$3 * y1\$2 * z1\$3 + \\
& x1\$2 * y1\$3 * z2\$3 + \\
& x1\$3 * y2\$3 * z1\$2 - \\
& x2\$3 * y1\$3 * z1\$2 - \\
& x1\$2 * y2\$3 * z1\$3 - \\
& x1\$3 * y1\$2 * z2\$3) / d, \\
& (x1\$1 * y2\$3 * z1\$3 + \\
& x2\$3 * y1\$3 * z1\$1 + \\
& x1\$3 * y1\$1 * z2\$3 - \\
& x1\$1 * y1\$3 * z2\$3 - \\
& x2\$3 * y1\$1 * z1\$3 - \\
& x1\$3 * y2\$3 * z1\$1) / d, \\
& (x1\$1 * y1\$2 * z2\$3 + \\
& x1\$2 * y2\$3 * z1\$1 + \\
& x2\$3 * y1\$1 * z1\$2 - \\
& x1\$1 * y2\$3 * z1\$2 - \\
& x1\$2 * y1\$1 * z2\$3 - \\
& x2\$3 * y1\$2 * z1\$1) / d]]))
\end{aligned}$$

lemma *set3-neqI*:

assumes $a \notin \{x, y, z\} \vee b \notin \{x, y, z\} \vee c \notin \{x, y, z\}$
shows $\{a, b, c\} \neq \{x, y, z\}$
using *assms* **by** *auto*

lemma *set4-neqI*:

assumes $a \notin \{w, x, y, z\} \vee b \notin \{w, x, y, z\} \vee c \notin \{w, x, y, z\} \vee d \notin \{w, x, y, z\}$

shows $\{a, b, c, d\} \neq \{w, x, y, z\}$
using *assms* **by** *auto*

lemma *orthogonal-matrix33*:

fixes $Q :: \text{real}^3 \times \text{real}^3$

shows (*orthogonal-matrix* Q) =

$(Q_{11} * Q_{11} + Q_{21} * Q_{21} + Q_{31} * Q_{31} = 1 \wedge$
 $Q_{11} * Q_{12} + Q_{21} * Q_{22} + Q_{31} * Q_{32} = 0 \wedge$
 $Q_{11} * Q_{13} + Q_{21} * Q_{23} + Q_{31} * Q_{33} = 0 \wedge$
 $Q_{12} * Q_{11} + Q_{22} * Q_{21} + Q_{32} * Q_{31} = 0 \wedge$
 $Q_{12} * Q_{12} + Q_{22} * Q_{22} + Q_{32} * Q_{32} = 1 \wedge$
 $Q_{12} * Q_{13} + Q_{22} * Q_{23} + Q_{32} * Q_{33} = 0 \wedge$
 $Q_{13} * Q_{11} + Q_{23} * Q_{21} + Q_{33} * Q_{31} = 0 \wedge$
 $Q_{13} * Q_{12} + Q_{23} * Q_{22} + Q_{33} * Q_{32} = 0 \wedge$
 $Q_{13} * Q_{13} + Q_{23} * Q_{23} + Q_{33} * Q_{33} = 1)$

by (*simp add: orthogonal-matrix transpose-def matrix-matrix-mult-def sum-3 mat-def*
vec-eq-iff forall-3)

lemma *matrix-by-cramer*:

fixes $A :: \text{real}^3 \times \text{real}^3$

assumes $A = \text{mk-matrix33 } x1 \ y1 \ z1 \ x2 \ y2 \ z2$

assumes $\det(\text{vector } [x1, y1, z1] :: \text{real}^3) \neq 0$

shows $A * v \ x1 = x2 \ A * v \ y1 = y2 \ A * v \ z1 = z2$

using *assms(1) matrix-by-cramer-lemma[OF assms(2)]*

unfolding *mk-matrix33-def Let-def*

by (*auto simp add: det-3 vec-eq-iff forall-3*)

lemma *matrix-by-cramer-orthog*:

fixes $x1 \ y1 \ z1 \ x2 \ y2 \ z2 :: \text{real}^3$

assumes *orthogonal-matrix* (*mk-matrix33* $x1 \ y1 \ z1 \ x2 \ y2 \ z2$)

assumes $\det(\text{vector } [x1, y1, z1] :: \text{real}^3) \neq 0$

shows *orthogonal-matrix* (*mk-matrix33* $x1 \ y1 \ z1 \ x2 \ y2 \ z2$) $\wedge$

$(*v) (\text{mk-matrix33 } x1 \ y1 \ z1 \ x2 \ y2 \ z2) \ ' \ \{x1, y1, z1\} = \{x2, y2, z2\}$

using *assms(1) matrix-by-cramer[OF - assms(2)]* **by** *auto*

Lookup table mapping a face to three vertices on that face that are adjacent (vertices that define a pair of intersecting edges)

definition *adjacency-table* :: $((\text{real}, 3) \text{ vec set} \times ((\text{real}, 3) \text{ vec} \times (\text{real}, 3) \text{ vec} \times (\text{real}, 3) \text{ vec})) \text{ list}$ **where**

adjacency-table = [
 $(\{\text{vector}[-1, -1, 1], \text{vector}[-1, 1, -1], \text{vector}[1, -1, -1]\}, (\text{vector}[-1, -1, 1], \text{vector}[-1, 1, -1], \text{vector}[1, -1, -1])),$
 $(\{\text{vector}[-1, -1, 1], \text{vector}[-1, 1, -1], \text{vector}[1, 1, 1]\}, (\text{vector}[-1, -1, 1], \text{vector}[-1, 1, -1], \text{vector}[1, 1, 1])),$
 $(\{\text{vector}[-1, -1, 1], \text{vector}[1, -1, -1], \text{vector}[1, 1, 1]\}, (\text{vector}[-1, -1, 1], \text{vector}[1, -1, -1], \text{vector}[1, 1, 1])),$
 $(\{\text{vector}[-1, 1, -1], \text{vector}[1, -1, -1], \text{vector}[1, 1, 1]\}, (\text{vector}[-1, 1, -1], \text{vector}[1, -1, -1], \text{vector}[1, 1, 1])),$
 $(\{\text{vector}[-1, 1, -1], \text{vector}[1, -1, -1], \text{vector}[1, 1, 1]\}, (\text{vector}[-1, 1, -1], \text{vector}[1, -1, -1], \text{vector}[1, 1, 1]))],$

$(\{vector[-1, -1, -1], vector[-1, -1, 1], vector[-1, 1, -1], vector[-1, 1, 1]\}, (vector[-1, -1, -1], vector[-1, -1, 1], vector[-1, 1, 1])),$
 $(\{vector[-1, -1, -1], vector[-1, -1, 1], vector[1, -1, -1], vector[1, -1, 1]\}, (vector[-1, -1, -1], vector[-1, -1, 1], vector[1, -1, 1])),$
 $(\{vector[-1, -1, -1], vector[-1, 1, -1], vector[1, -1, -1], vector[1, 1, -1]\}, (vector[-1, -1, -1], vector[-1, 1, -1], vector[1, 1, -1])),$
 $(\{vector[-1, -1, 1], vector[-1, 1, 1], vector[1, -1, 1], vector[1, 1, 1]\}, (vector[-1, -1, 1], vector[-1, 1, 1], vector[1, 1, 1])),$
 $(\{vector[-1, 1, -1], vector[-1, 1, 1], vector[1, 1, -1], vector[1, 1, 1]\}, (vector[-1, 1, -1], vector[-1, 1, 1], vector[1, 1, 1])),$
 $(\{vector[1, -1, -1], vector[1, -1, 1], vector[1, 1, -1], vector[1, 1, 1]\}, (vector[1, -1, -1], vector[1, -1, 1], vector[1, 1, 1])),$

$(\{vector[-1, 0, 0], vector[0, -1, 0], vector[0, 0, -1]\}, (vector[-1, 0, 0], vector[0, -1, 0], vector[0, 0, -1])),$
 $(\{vector[-1, 0, 0], vector[0, -1, 0], vector[0, 0, 1]\}, (vector[-1, 0, 0], vector[0, -1, 0], vector[0, 0, 1])),$
 $(\{vector[-1, 0, 0], vector[0, 1, 0], vector[0, 0, -1]\}, (vector[-1, 0, 0], vector[0, 1, 0], vector[0, 0, -1])),$
 $(\{vector[-1, 0, 0], vector[0, 1, 0], vector[0, 0, 1]\}, (vector[-1, 0, 0], vector[0, 1, 0], vector[0, 0, 1])),$
 $(\{vector[1, 0, 0], vector[0, -1, 0], vector[0, 0, -1]\}, (vector[1, 0, 0], vector[0, -1, 0], vector[0, 0, -1])),$
 $(\{vector[1, 0, 0], vector[0, -1, 0], vector[0, 0, 1]\}, (vector[1, 0, 0], vector[0, -1, 0], vector[0, 0, 1])),$
 $(\{vector[1, 0, 0], vector[0, 1, 0], vector[0, 0, -1]\}, (vector[1, 0, 0], vector[0, 1, 0], vector[0, 0, -1])),$
 $(\{vector[1, 0, 0], vector[0, 1, 0], vector[0, 0, 1]\}, (vector[1, 0, 0], vector[0, 1, 0], vector[0, 0, 1])),$

$(\{vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)], vector[rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 (-1) 0, rat5 (-1) 0, rat5 (-1) 0], vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0]\}, (vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)], vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0])),$

$(\{vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)], vector[rat5 (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0], vector[rat5 (-1) 0, rat5 1 0, rat5 (-1) 0], vector[rat5 (-1) 0, rat5 1 0, rat5 1 0]\}, (vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)], vector[rat5 (-1) 0, rat5 1 0, rat5 1 0])),$

$(\{vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0], vector[rat5 (-1) 0, rat5 1 0, rat5 1 0], vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (1/2) (1/2)], vector[rat5 0 0, rat5 (1/2) (-1/2), rat5 (1/2) (1/2)]\}, (vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0], vector[rat5 (-1) 0, rat5 1 0, rat5 1 0])),$

[illegible]

[illegible]

```

vector[rat5 0 0, rat5 1 0, rat5 (-1/2) (-1/2)])),
  ({vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1) 0], vector[rat5 0 0, rat5
(-1) 0, rat5 (-1/2) (-1/2)], vector[rat5 0 0, rat5 1 0, rat5 (-1/2) (-1/2)]},
(vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1) 0], vector[rat5 0 0, rat5 (-1) 0,
rat5 (-1/2) (-1/2)], vector[rat5 0 0, rat5 1 0, rat5 (-1/2) (-1/2)])),
  ({vector[rat5 (1/2) (1/2), rat5 0 0, rat5 1 0], vector[rat5 1 0, rat5 (-1/2)
(-1/2), rat5 0 0], vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2) (1/2)]}, (vector[rat5
(1/2) (1/2), rat5 0 0, rat5 1 0], vector[rat5 1 0, rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2) (1/2)])),
  ({vector[rat5 (1/2) (1/2), rat5 0 0, rat5 1 0], vector[rat5 1 0, rat5 (1/2)
(1/2), rat5 0 0], vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)]}, (vector[rat5 (1/2)
(1/2), rat5 0 0, rat5 1 0], vector[rat5 1 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5
0 0, rat5 1 0, rat5 (1/2) (1/2)])),
  ({vector[rat5 (1/2) (1/2), rat5 0 0, rat5 1 0], vector[rat5 0 0, rat5 (-1)
0, rat5 (1/2) (1/2)], vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)]}, (vector[rat5
(1/2) (1/2), rat5 0 0, rat5 1 0], vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2) (1/2)],
vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)])),
  ({vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 1 0, rat5
(-1/2) (-1/2), rat5 0 0], vector[rat5 0 0, rat5 (-1) 0, rat5 (-1/2) (-1/2)]},
(vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 1 0, rat5 (-1/2)
(-1/2), rat5 0 0], vector[rat5 0 0, rat5 (-1) 0, rat5 (-1/2) (-1/2)])),
  ({vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 1 0, rat5
(-1/2) (-1/2), rat5 0 0], vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2) (1/2)]},
(vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 1 0, rat5 (-1/2)
(-1/2), rat5 0 0], vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2) (1/2)])),
  ({vector[rat5 (-1) 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 1 0, rat5 (1/2)
(1/2), rat5 0 0], vector[rat5 0 0, rat5 1 0, rat5 (-1/2) (-1/2)]}, (vector[rat5
(-1) 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 1 0, rat5 (1/2) (1/2), rat5 0 0],
vector[rat5 0 0, rat5 1 0, rat5 (-1/2) (-1/2)])),
  ({vector[rat5 (-1) 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 1 0, rat5 (1/2)
(1/2), rat5 0 0], vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)]}, (vector[rat5 (-1)
0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 1 0, rat5 (1/2) (1/2), rat5 0 0],
vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)]))
]

```

definition *adjacent-points* :: (real, 3) vec set $\Rightarrow$ (real, 3) vec $\times$ (real, 3) vec $\times$ (real, 3) vec **where**

```

adjacent-points S = (
  case map-of adjacency-table S of
    Some res  $\Rightarrow$  res
  | None  $\Rightarrow$  undefined
)

```

lemma *set43-neqI*:

```

assumes a  $\notin \{x, y, z\} \vee b \notin \{x, y, z\} \vee c \notin \{x, y, z\} \vee d \notin \{x, y, z\}$ 
shows {a, b, c, d}  $\neq \{x, y, z\}$ 
      {x, y, z}  $\neq \{a, b, c, d\}$ 
using assms by auto

```

lemma *solids-vectors-to-rat5*:

shows $\text{vector}[1, 1, 1] = \text{vector}[\text{rat5 } (1) \ 0, \text{rat5 } (1) \ 0, \text{rat5 } (1) \ 0]$
 $\text{vector}[1, 1, -1] = \text{vector}[\text{rat5 } (1) \ 0, \text{rat5 } (1) \ 0, \text{rat5 } (-1) \ 0]$
 $\text{vector}[1, -1, 1] = \text{vector}[\text{rat5 } (1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1) \ 0]$
 $\text{vector}[1, -1, -1] = \text{vector}[\text{rat5 } (1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0]$
 $\text{vector}[-1, 1, 1] = \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (1) \ 0, \text{rat5 } (1) \ 0]$
 $\text{vector}[-1, 1, -1] = \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (1) \ 0, \text{rat5 } (-1) \ 0]$
 $\text{vector}[-1, -1, 1] = \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1) \ 0]$
 $\text{vector}[-1, -1, -1] = \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0]$
 $\text{vector}[1, 0, 0] = \text{vector}[\text{rat5 } (1) \ 0, \text{rat5 } (0) \ 0, \text{rat5 } (0) \ 0]$
 $\text{vector}[-1, 0, 0] = \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (0) \ 0, \text{rat5 } (0) \ 0]$
 $\text{vector}[0, 0, 1] = \text{vector}[\text{rat5 } (0) \ 0, \text{rat5 } (0) \ 0, \text{rat5 } (1) \ 0]$
 $\text{vector}[0, 0, -1] = \text{vector}[\text{rat5 } (0) \ 0, \text{rat5 } (0) \ 0, \text{rat5 } (-1) \ 0]$
 $\text{vector}[0, 1, 0] = \text{vector}[\text{rat5 } (0) \ 0, \text{rat5 } (1) \ 0, \text{rat5 } (0) \ 0]$
 $\text{vector}[0, -1, 0] = \text{vector}[\text{rat5 } (0) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (0) \ 0]$
unfolding *rat5-def*
by *auto*

lemma *set53-neqI*:

assumes $a \notin \{x, y, z\} \vee b \notin \{x, y, z\} \vee c \notin \{x, y, z\} \vee d \notin \{x, y, z\} \vee e \notin \{x, y, z\}$
shows $\{a, b, c, d, e\} \neq \{x, y, z\}$
 $\{x, y, z\} \neq \{a, b, c, d, e\}$
using *assms* **by** *auto*

lemma *set54-neqI*:

assumes $a \notin \{w, x, y, z\} \vee b \notin \{w, x, y, z\} \vee c \notin \{w, x, y, z\} \vee d \notin \{w, x, y, z\} \vee e \notin \{w, x, y, z\}$
shows $\{a, b, c, d, e\} \neq \{w, x, y, z\}$
 $\{w, x, y, z\} \neq \{a, b, c, d, e\}$
using *assms* **by** *auto*

lemma *set5-neqI*:

assumes $a \notin \{v, w, x, y, z\} \vee b \notin \{v, w, x, y, z\} \vee c \notin \{v, w, x, y, z\} \vee d \notin \{v, w, x, y, z\} \vee e \notin \{v, w, x, y, z\}$
shows $\{a, b, c, d, e\} \neq \{v, w, x, y, z\}$
using *assms* **by** *auto*

lemma *set3-eq-iff*:

shows $(\{a, b, c\} = \{d, e, f\}) =$
 $(a \in \{d, e, f\} \wedge b \in \{d, e, f\} \wedge c \in \{d, e, f\} \wedge$
 $d \in \{a, b, c\} \wedge e \in \{a, b, c\} \wedge f \in \{a, b, c\})$
by *auto*

lemma *set3-vector-neq*:

assumes $(a1 \neq d1 \wedge a1 \neq d2 \wedge a1 \neq d3)$
shows $\{(\text{vector } [a1, b1, c1]::\text{real}^3), \text{vector } [a2, b2, c2], \text{vector } [a3, b3, c3]\} \neq$
 $\{\text{vector } [d1, e1, f1], \text{vector } [d2, e2, f2], \text{vector } [d3, e3, f3]\}$
by $(\text{auto simp add: set3-neqI assms vector3-neq})$

[illegible]

```

1 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)]}
= (vector[rat5 (-1) 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 1 0, rat5 (1/2)
(1/2), rat5 0 0], vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)])
apply (simp-all add: adjacent-points-def adjacency-table-def set3-neqI set4-neqI
set43-neqI
      set53-neqI set54-neqI set5-neqI vector3-neq)[18]
by (simp-all add: rat5-eq adjacent-points-def adjacency-table-def set3-neqI set4-neqI
set43-neqI
      set53-neqI set54-neqI set5-neqI vector3-neq set3-vector-neq solids-vectors-to-rat5)

```

definition *mk-matrix33-ap* :: $(\text{real}^3) \text{ set} \Rightarrow (\text{real}^3) \text{ set} \Rightarrow \text{real}^3 \times \text{real}^3$ **where**
mk-matrix33-ap f1 f2 = (let vs = adjacent-points f1; us = adjacent-points f2 in
 mk-matrix33 (fst vs) (fst (snd vs)) (snd (snd vs))
 (fst us) (fst (snd us)) (snd (snd us)))

```

method show-congruence =
  (((erule forw-subst)+)?;
  (match conclusion in convex hull (s::( $\text{real}^3$ ) set) congruent convex hull (t::( $\text{real}^3$ )
set) for s t  $\Rightarrow$ 
    <(rule congruent-simple), (rule exI[where x=mk-matrix33-ap s t])>);
  (simp only: mk-matrix33-ap-def Let-def fst-conv snd-conv adjacent-points);
  (rule matrix-by-cramer-orthog | simp add: matrix-by-cramer det-3 rat5-eq-0 rat5-add
rat5-mul rat5-sub);
  (simp only: mk-matrix33-def vector-3 det-3 rat5-add rat5-mul rat5-sub rat5-eq);
  (simp add: matrix-vector-mul-3 orthogonal-matrix33 rat5-add rat5-mul rat5-sub
rat5-div rat5-eq
    rat5-eq-0 rat5-eq-1 Let-def; fast))

```

4.7 Equiangularity

vangle is from Manuel Eberl's AFP entry on triangles (see <https://www.isa-afp.org/thys/Triangle/Angles.html>)

definition *vangle* :: '*a*::*real-inner* $\Rightarrow$ '*a* $\Rightarrow$ *real* **where**
vangle x y = (if x = 0 $\vee$ y = 0 then pi / 2 else arccos (x • y / (norm x * norm y)))

lemma *vangle-commute*:
shows *vangle* x y = *vangle* y x
unfolding *vangle-def* inner-commute[of y x]
by argo

lemma *isometry-maintains-angles-forw*:
fixes c
assumes orthogonal-transformation h
defines g $\equiv$ ($\lambda x. c + h x$)
shows ((e1 face-of f1 $\wedge$ aff-dim e1 = 1) $\wedge$
 (e2 face-of f1 $\wedge$ aff-dim e2 = 1) $\wedge$
 e1 = convex hull {v, a} $\wedge$ e2 = convex hull {v, b} $\wedge$
 vangle (a - v) (b - v) = t) $\implies$

$((g \text{ ' } e1 \text{ face-of } g \text{ ' } f1 \wedge \text{aff-dim } (g \text{ ' } e1) = 1) \wedge$
 $(g \text{ ' } e2 \text{ face-of } g \text{ ' } f1 \wedge \text{aff-dim } (g \text{ ' } e2) = 1) \wedge$
 $g \text{ ' } e1 = \text{convex hull } \{g \text{ v, } g \text{ a}\} \wedge g \text{ ' } e2 = \text{convex hull } \{g \text{ v, } g \text{ b}\} \wedge$
 $\text{vangle } (g \text{ a} - g \text{ v}) (g \text{ b} - g \text{ v}) = t)$
 (is ?lhs $\implies$?rhs)
proof –
 assume *: ?lhs
 then have $g \text{ ' } e1 \text{ face-of } g \text{ ' } f1 \wedge \text{aff-dim } (g \text{ ' } e1) = 1$
 $g \text{ ' } e2 \text{ face-of } g \text{ ' } f1 \wedge \text{aff-dim } (g \text{ ' } e2) = 1$
 using *face-of-isometry-eq*[*OF assms*(1)] *aff-dim-isometry-eq*[*OF assms*(1)] *g-def*
 by *auto*
 moreover have $g \text{ ' } e1 = \text{convex hull } \{g \text{ v, } g \text{ a}\}$
 $g \text{ ' } e2 = \text{convex hull } \{g \text{ v, } g \text{ b}\}$
 using * *convex-hull-isometry*[*OF assms*(1), of *c*] *g-def*
 by *simp-all*
 moreover have $\text{vangle } (g \text{ a} - g \text{ v}) (g \text{ b} - g \text{ v}) = t$
proof –
 have $a - v \neq 0 \text{ b} - v \neq 0$
 using * by *auto*
 have $\text{vangle } (g \text{ a} - g \text{ v}) (g \text{ b} - g \text{ v}) =$
 $\arccos ((g \text{ a} - g \text{ v}) \cdot (g \text{ b} - g \text{ v}) / (\text{norm } (g \text{ a} - g \text{ v}) * \text{norm } (g \text{ b} - g$
 $v)))$
 unfolding *vangle-def* by *simp*
 also have $\dots = \arccos ((h \text{ a} - h \text{ v}) \cdot (h \text{ b} - h \text{ v}) / (\text{norm } (h \text{ a} - h \text{ v}) * \text{norm}$
 $(h \text{ b} - h \text{ v})))$
 unfolding *g-def* by *simp*
 also have $\dots = \arccos (h (a - v) \cdot h (b - v) / (\text{norm } (h (a - v)) * \text{norm } (h$
 $(b - v))))$
 using *orthogonal-transformation-linear*[*OF assms*(1)]
 by (*simp add: linear-diff*)
 also have $\dots = \arccos ((a - v) \cdot (b - v) / (\text{norm } (a - v) * \text{norm } (b - v)))$
 unfolding *orthogonal-transformation-norm*[*OF assms*(1)]
 by (*metis* (*mono-tags*, *opaque-lifting*) *assms*(1) *orthogonal-transformation-def*)
 finally show ?thesis
 using * *vangle-def* $\langle a - v \neq 0 \rangle \langle b - v \neq 0 \rangle$
 by *metis*
qed
 ultimately show ?rhs
 by *blast*
qed

definition *equiangular* :: '*a::euclidean-space set* $\Rightarrow$ *real* $\Rightarrow$ *bool*
 where *equiangular* *f t* $\longleftrightarrow$
 $(\forall e1 \text{ e2 } v. (e1 \text{ face-of } f \wedge \text{aff-dim } e1 = 1) \wedge (e2 \text{ face-of } f \wedge \text{aff-dim } e2 =$
 $1) \wedge$
 $e1 \neq e2 \wedge v \text{ extreme-point-of } e1 \wedge v \text{ extreme-point-of } e2 \longrightarrow$
 $(\exists a \text{ b. } e1 = \text{convex hull } \{v, a\} \wedge e2 = \text{convex hull } \{v, b\} \wedge$
 $\text{vangle } (a - v) (b - v) = t))$

```

lemma equiangular-congruent-forw:
  assumes f1 congruent f2
  assumes equiangular f1 t
  shows equiangular f2 t
proof -
  {
    fix e1' e2' v'
    assume *: (e1' face-of f2  $\wedge$  aff-dim e1' = 1)  $\wedge$ 
      (e2' face-of f2  $\wedge$  aff-dim e2' = 1)  $\wedge$ 
      e1'  $\neq$  e2'  $\wedge$  v' extreme-point-of e1'  $\wedge$  v' extreme-point-of e2'
    obtain c h where orthog:orthogonal-transformation h and f2: f2 = ( $\lambda$ x. c +
h x) ' f1
    using assms unfolding congruent-def by blast
    then have orthog-inv: orthogonal-transformation (inv h)
    using orthogonal-transformation-inv by blast
    have bounded-linear (inv h)
    using orthogonal-transformation-linear[OF orthog-inv]
    by (simp add: linear-conv-bounded-linear)
    have ( $\lambda$ y. - inv h c + inv h y) ' f2 = ( $\lambda$ x. - inv h c + inv h (c + h x)) ' f1
    using f2 by blast
    also have ... = ( $\lambda$ x. - inv h c + inv h c + inv h (h x)) ' f1
    using  $\langle$ bounded-linear (inv h) $\rangle$ 
    by (simp add: linear-simps)
    also have ... = ( $\lambda$ x. inv h (h x)) ' f1
    by simp
    finally have f1: f1 = ( $\lambda$ y. - inv h c + inv h y) ' f2
    by (simp add: orthog bijection.intro bijection.inv-left orthogonal-transformation-bij)

    define g where g = ( $\lambda$ x. c + h x)
    define g-inv where g-inv = ( $\lambda$ x. - inv h c + inv h x)
    have g-inv-g: g-inv (g y) = y for y
    using g-def g-inv-def
    by (simp add:  $\langle$ bounded-linear (inv h) $\rangle$  linear-simps(1) orthog orthog-
nal-transformation-inj)
    then have g-g-inv: g (g-inv x) = x for x
    using g-def g-inv-def
    by (metis add-minus-cancel bijection.intro
      bijection.inv-right orthog orthog-transformation-bij)

    define e1 where e1 = g-inv ' e1'
    define e2 where e2 = g-inv ' e2'
    define v where v = g-inv v'

    have e1-e2: e1 face-of f1  $\wedge$  aff-dim e1 = 1 e2 face-of f1  $\wedge$  aff-dim e2 = 1
    using * e1-def e2-def f1 face-of-isometry-eq[OF orthog-inv]
      aff-dim-isometry-eq[OF orthog-inv]
    unfolding g-def g-inv-def
    by presburger+
    moreover have e1  $\neq$  e2

```

```

    using e1-def e2-def * g-inv-g g-g-inv
    by (metis (no-types) image-f-inv-f surj-def surj-imp-inv-eq)
  moreover have v extreme-point-of e1 v extreme-point-of e2
    unfolding e1-def e2-def v-def using * extreme-point-of-isometry-eq[OF orthog-inv] g-inv-def
    by blast+
  ultimately obtain a b where ab: e1 = convex hull {v, a} e2 = convex hull {v, b}
    vangle (a - v) (b - v) = t
    using assms(2) unfolding equiangular-def
    by meson

  then have g ' e1 = convex hull {g v, g a} ∧ g ' e2 = convex hull {g v, g b} ∧
    vangle (g a - g v) (g b - g v) = t
    using isometry-maintains-angles-forw[OF orthog, of e1 f1 e2 v] g-def e1-e2
    by presburger
  then have ∃ a b. e1' = convex hull {v', a} ∧ e2' = convex hull {v', b} ∧
    vangle (a - v') (b - v') = t
    using * unfolding e1-def e2-def v-def
    by (metis (mono-tags, lifting) image-f-inv-f g-inv-g g-g-inv surj-def surj-imp-inv-eq)
} then show ?thesis
  unfolding equiangular-def
  by blast
qed

```

```

lemma equiangular-congruent:
  assumes f1 congruent f2
  shows equiangular f1 t ⟷ equiangular f2 t
  using equiangular-congruent-forw assms congruent-sym
  by blast

```

end

5 Tetrahedron

```

theory Tetrahedron
  imports Computation
begin

```

5.1 Definition

```

definition std-tetrahedron :: (real^3) set
  where std-tetrahedron ≡ convex hull
    {vector [1, 1, 1], vector [-1, -1, 1],
     vector [-1, 1, -1], vector [1, -1, -1]}

```

```

lemma tetrahedron-fulldim:
  shows aff-dim std-tetrahedron = 3
  unfolding std-tetrahedron-def

```

apply (*rule polytope-3D*)
by (*simp add: vector3-sub vector3-cross vector3-dot vector3-eq-0*)**+**

lemma *tetrahedron-polyhedron*:
shows *polyhedron std-tetrahedron*
by (*simp add: std-tetrahedron-def polytope-convex-hull polytope-imp-polyhedron*)

5.2 Facets, edges, and vertices

lemma *tetrahedron-facets*:
 $f \text{ face-of } \text{std-tetrahedron} \wedge \text{aff-dim } f = 2 \longleftrightarrow$
 $f = \text{convex hull } \{\text{vector}[-1, -1, 1], \text{vector}[-1, 1, -1], \text{vector}[1, -1, -1]\}$
 $\vee$
 $f = \text{convex hull } \{\text{vector}[-1, -1, 1], \text{vector}[-1, 1, -1], \text{vector}[1, 1, 1]\} \vee$
 $f = \text{convex hull } \{\text{vector}[-1, -1, 1], \text{vector}[1, -1, -1], \text{vector}[1, 1, 1]\} \vee$
 $f = \text{convex hull } \{\text{vector}[-1, 1, -1], \text{vector}[1, -1, -1], \text{vector}[1, 1, 1]\}$
unfolding *std-tetrahedron-def*
apply (*simp only: compute-faces-2 finite-insert finite.emptyI*)
apply (*simp only: compute-faces-2-step-1 bex-empty simp-thms(31)*)
apply (*simp only: compute-faces-2-step-2 bex-empty simp-thms(31)*)
apply (*simp only: empty-iff ball-insert bex-empty bex-simps(5) ball-simps(5)*)
simp-thms(21,31)
apply (*simp add: vector3-sub vector3-cross vector3-eq-0 vector3-dot Let-def*)
apply (*simp only: insert-commute*)
by *linarith*

lemma *tetrahedron-facet1-edges*:
defines $v1 \equiv \text{vector}[-1, -1, 1] :: \text{real}^3$
and $v2 \equiv \text{vector}[-1, 1, -1] :: \text{real}^3$
and $v3 \equiv \text{vector}[1, -1, -1] :: \text{real}^3$
and $n \equiv \text{vector}[-4, -4, -4] :: \text{real}^3$
shows ($e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1$) =
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
proof –
have $n \cdot n \neq 0$
unfolding *n-def* **by** (*simp add: vector3-eq-0*)
have *plane*: $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = 4$
by (*auto simp add: assms vector3-dot*)
have *finite*: $\text{finite } \{v1, v2, v3\}$
using *assms* **by** *blast*
show *?thesis*
by (*facet-edges plane: plane n-nz: n-nz defs: assms*)
qed

lemma *tetrahedron-facet2-edges*:
defines $v1 \equiv \text{vector}[-1, -1, 1] :: \text{real}^3$
and $v2 \equiv \text{vector}[-1, 1, -1] :: \text{real}^3$
and $v3 \equiv \text{vector}[1, 1, 1] :: \text{real}^3$

```

    and  $n \equiv \text{vector}[-4, 4, 4] :: \text{real}^3$ 
    shows  $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$ 
       $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$ 
  proof -
    have  $n\text{-nz}: n \neq 0$ 
    unfolding  $n\text{-def}$  by (simp add: vector3-eq-0)
    have  $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = 4$ 
    by (auto simp add: assms vector3-dot)
    have  $\text{finite}: \text{finite } \{v1, v2, v3\}$ 
    using assms by blast
    show ?thesis
    by (facet-edges plane: plane  $n\text{-nz}$ :  $n\text{-nz}$  defs: assms)
  qed

```

```

lemma tetrahedron-facet3-edges:
  defines  $v1 \equiv \text{vector}[-1, -1, 1] :: \text{real}^3$ 
    and  $v2 \equiv \text{vector}[1, -1, -1] :: \text{real}^3$ 
    and  $v3 \equiv \text{vector}[1, 1, 1] :: \text{real}^3$ 
    and  $n \equiv \text{vector}[4, -4, 4] :: \text{real}^3$ 
  shows  $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$ 
     $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$ 
  proof -
    have  $n\text{-nz}: n \neq 0$ 
    unfolding  $n\text{-def}$  by (simp add: vector3-eq-0)
    have  $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = 4$ 
    by (auto simp add: assms vector3-dot)
    have  $\text{finite}: \text{finite } \{v1, v2, v3\}$ 
    using assms by blast
    show ?thesis
    by (facet-edges plane: plane  $n\text{-nz}$ :  $n\text{-nz}$  defs: assms)
  qed

```

```

lemma tetrahedron-facet4-edges:
  defines  $v1 \equiv \text{vector}[-1, 1, -1] :: \text{real}^3$ 
    and  $v2 \equiv \text{vector}[1, -1, -1] :: \text{real}^3$ 
    and  $v3 \equiv \text{vector}[1, 1, 1] :: \text{real}^3$ 
    and  $n \equiv \text{vector}[4, 4, -4] :: \text{real}^3$ 
  shows  $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$ 
     $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$ 
  proof -
    have  $n\text{-nz}: n \neq 0$ 
    unfolding  $n\text{-def}$  by (simp add: vector3-eq-0)
    have  $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = 4$ 
    by (auto simp add: assms vector3-dot)
    have  $\text{finite}: \text{finite } \{v1, v2, v3\}$ 
    using assms by blast
  end

```

```

show ?thesis
  by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemmas tetrahedron-facet-edges =
  tetrahedron-facet1-edges tetrahedron-facet2-edges
  tetrahedron-facet3-edges tetrahedron-facet4-edges

lemma tetrahedron-edges-per-face:
  assumes f face-of std-tetrahedron
  and aff-dim f = 2
  shows card {e. e face-of std-tetrahedron ∧ aff-dim e = 1 ∧ e ⊆ f} = 3
  using conjI[OF assms] unfolding tetrahedron-facets face-edge-set-eq[OF assms(1)]
  apply (elim disjE forw-subst)
  unfolding tetrahedron-facet-edges set-cases card-3-iff2
    segment-convex-hull[symmetric] closed-segment-eq
  by (simp-all add: vector3-eq doubleton-eq-iff)

lemma tetrahedron-edges:
  e face-of std-tetrahedron ∧ aff-dim e = 1 ⟷
  e = convex hull {vector [-1, -1, 1], vector [-1, 1, -1]} ∨
  e = convex hull {vector [-1, -1, 1], vector [1, -1, -1]} ∨
  e = convex hull {vector [-1, -1, 1], vector [1, 1, 1]} ∨
  e = convex hull {vector [-1, 1, -1], vector [1, -1, -1]} ∨
  e = convex hull {vector [-1, 1, -1], vector [1, 1, 1]} ∨
  e = convex hull {vector [1, -1, -1], vector [1, 1, 1]}
  (is ?lhs ⟷ ?rhs)
proof
  assume *: ?lhs
  then obtain f where f-facts: f face-of std-tetrahedron ∧ aff-dim f = 2 e face-of
  f ∧ aff-dim e = 1
  using tetrahedron-polyhedron tetrahedron-fulldim edge-belongs-to-face
  by blast
  show ?rhs
  using f-facts unfolding tetrahedron-facets
  apply (elim disjE) apply (all ‹hypsubst-thin›)
  unfolding tetrahedron-facet-edges
  by linarith+
next
  assume ?rhs
  then show ?lhs
  using tetrahedron-facets tetrahedron-facet-edges
  by (smt (verit, ccfv-threshold) face-of-trans insertI1 insert-commute mem-Collect-eq)
qed

lemma tetrahedron-vertices:
  shows (v face-of std-tetrahedron ∧ aff-dim v = 0) =
    (v = {vector [-1, -1, 1]} ∨
     v = {vector [-1, 1, -1]} ∨
     v = {vector [1, -1, -1]} ∨
     v = {vector [1, 1, 1]})

```

```

      v = {vector [1, - 1, - 1]} ∨
      v = {vector [1, 1, 1]}
    (is ?lhs = ?rhs)
  proof
    assume *: ?lhs
    then have aff-dim v = 0
      by blast
    obtain e where e-facts: e face-of std-tetrahedron ∧ aff-dim e = 1 v face-of e
      using * tetrahedron-polyhedron tetrahedron-fulldim vertex-belongs-to-edge
      by blast
    show ?rhs
      using e-facts(1)
      unfolding tetrahedron-edges
      using vertices-of-edge[OF - ⟨v face-of e⟩ ⟨aff-dim v = 0⟩]
      by metis
  next
    assume ?rhs
    then show ?lhs
      apply (elim disjE)
      by (metis (no-types, opaque-lifting) aff-dim-sing extreme-point-of-segment
        face-of-singleton face-of-trans segment-convex-hull tetrahedron-edges)+
  qed

lemma tetrahedron-edges-per-vertex:
  assumes v face-of std-tetrahedron
    and aff-dim v = 0
  shows card {e. e face-of std-tetrahedron ∧ aff-dim e = 1 ∧ v ⊆ e} = 3
  unfolding vertex-edge-set-eq[OF assms(1)]
  using conjI[OF assms]

  unfolding tetrahedron-vertices conj-assoc[symmetric] tetrahedron-edges
  by (elim disjE forw-subst) edges-per-vertex+

```

5.3 Regularity

```

lemma tetrahedron-congruent-edges:
  assumes e1 face-of std-tetrahedron ∧ aff-dim e1 = 1
    and e2 face-of std-tetrahedron ∧ aff-dim e2 = 1
  shows e1 congruent e2
proof -
  let ?edges =
    {convex hull {(vector [-1, -1, 1]::real^3), vector [-1, 1, -1]},
     convex hull {vector [-1, -1, 1], vector [1, -1, -1]},
     convex hull {vector [-1, -1, 1], vector [1, 1, 1]},
     convex hull {vector [-1, 1, -1], vector [1, -1, -1]},
     convex hull {vector [-1, 1, -1], vector [1, 1, 1]},
     convex hull {vector [1, -1, -1], vector [1, 1, 1]}}
  have e1 ∈ ?edges e2 ∈ ?edges
    using assms unfolding tetrahedron-edges

```

```

    by auto
    moreover have  $\forall e \in ?edges. \text{convex hull } \{\text{vector } [-1, -1, 1], \text{vector } [-1, 1, -1]\} \text{ congruent } e$ 
    apply (simp only: ball-simps(7) Set.ball-empty simp-thms(21))
    unfolding conj-assoc[symmetric] segment-convex-hull[symmetric]
    apply (rule conjI)+
    apply (all <rule congruent-segments>)
    unfolding dist-norm vector3-sub
    apply (simp-all add: norm-eq)
    unfolding vector3-dot
    by linarith+
    ultimately show ?thesis
    using congruent-set by meson
qed

```

lemma *tetrahedron-congruent-faces:*

```

    assumes  $f1 \text{ face-of std-tetrahedron} \wedge \text{aff-dim } f1 = 2$ 
    and  $f2 \text{ face-of std-tetrahedron} \wedge \text{aff-dim } f2 = 2$ 
    shows  $f1 \text{ congruent } f2$ 
proof -
    let ?faces =
    {convex hull  $\{(\text{vector}[-1, -1, 1]::\text{real}^3), \text{vector}[-1, 1, -1], \text{vector}[1, -1, -1]\}$ ,
    convex hull  $\{\text{vector}[-1, -1, 1], \text{vector}[-1, 1, -1], \text{vector}[1, 1, 1]\}$ ,
    convex hull  $\{\text{vector}[-1, -1, 1], \text{vector}[1, -1, -1], \text{vector}[1, 1, 1]\}$ ,
    convex hull  $\{\text{vector}[-1, 1, -1], \text{vector}[1, -1, -1], \text{vector}[1, 1, 1]\}$ }
    have  $f1 \in ?faces \wedge f2 \in ?faces$ 
    using assms tetrahedron-facets
    by auto
    moreover have  $\forall f \in ?faces. \text{convex hull } \{\text{vector}[-1, -1, 1], \text{vector}[-1, 1, -1], \text{vector}[1, -1, -1]\} \text{ congruent } f$ 
    apply (simp only: ball-simps(7) Set.ball-empty simp-thms(21))
    unfolding conj-assoc[symmetric]
    by (rule conjI)+ show-congruence+
    ultimately show ?thesis
    using congruent-set by meson
qed

```

lemma *tetrahedron-facet1-equiangular:*

```

    defines  $v1 \equiv \text{vector}[-1, -1, 1] :: \text{real}^3$ 
    defines  $v2 \equiv \text{vector}[-1, 1, -1] :: \text{real}^3$ 
    defines  $v3 \equiv \text{vector}[1, -1, -1] :: \text{real}^3$ 
    shows  $\text{equiangular } (\text{convex hull } \{v1, v2, v3\}) (\pi / 3)$ 
    unfolding equiangular-def
proof (clarify)
    fix  $e1 \ e2 \ v$ 
    assume *:  $e1 \text{ face-of convex hull } \{v1, v2, v3\} \wedge \text{aff-dim } e1 = 1$ 
    and  $e2 \text{ face-of convex hull } \{v1, v2, v3\} \wedge \text{aff-dim } e2 = 1$ 

```

```

      e1 ≠ e2 v extreme-point-of e1 v extreme-point-of e2
have (v = v1 ∧ ((e1 = convex hull {v1, v2} ∧ e2 = convex hull {v1, v3}) ∨
      (e1 = convex hull {v1, v3} ∧ e2 = convex hull {v1, v2}))) ∨
      (v = v2 ∧ ((e1 = convex hull {v2, v1} ∧ e2 = convex hull {v2, v3}) ∨
      (e1 = convex hull {v2, v3} ∧ e2 = convex hull {v2, v1}))) ∨
      (v = v3 ∧ ((e1 = convex hull {v3, v1} ∧ e2 = convex hull {v3, v2}) ∨
      (e1 = convex hull {v3, v2} ∧ e2 = convex hull {v3, v1})))
using conjI[OF *(1,2)] conjI[OF *(3,4)] *(5-7) unfolding assms tetrahe-
dron-facet1-edges
by (smt (verit, best) extreme-point-of-convex-hull-2 insert-commute vector3-eq)
moreover have vangle (v2 - v1) (v3 - v1) = pi / 3
      vangle (v1 - v2) (v3 - v2) = pi / 3
      vangle (v1 - v3) (v2 - v3) = pi / 3
unfolding assms vector3-sub norm-eq-sqrt-inner vector3-dot vangle-def
by (simp-all add: vector3-eq-0)
ultimately show (∃ a b. e1 = convex hull {v, a} ∧ e2 = convex hull {v, b} ∧
      vangle (a - v) (b - v) = pi / 3)
using vangle-commute
by metis
qed

```

lemma tetrahedron-equangular:

```

assumes f face-of std-tetrahedron ∧ aff-dim f = 2
shows equiangular f (pi / 3)
proof -
  define f1 where f1 = convex hull {vector[-1, -1, 1]::real^3, vector[-1, 1,
-1], vector[1, -1, -1]}
  have f1 face-of std-tetrahedron ∧ aff-dim f1 = 2
    using f1-def tetrahedron-facets by blast
  then have f1 congruent f
    using assms tetrahedron-congruent-faces by blast
  then show ?thesis
    using equiangular-congruent tetrahedron-facet1-equiangular f1-def
    by blast
qed

```

end

6 Cube

```

theory Cube
imports Computation
begin

```

6.1 Definition

```

definition std-cube :: (real, 3) vec set
where std-cube ≡ convex hull
  {vector [1, 1, 1], vector [1, 1, -1],

```

```

vector [1, -1, 1], vector [1, -1, -1],
vector [-1, 1, 1], vector [-1, 1, -1],
vector [-1, -1, 1], vector [-1, -1, -1]}

```

lemma *cube-fulldim*:

shows *aff-dim std-cube = 3*

unfolding *std-cube-def*

apply (*rule polytope-3D*)

by (*simp add: vector3-sub vector3-cross vector3-dot vector3-eq-0*)**+**

lemma *cube-polyhedron*:

shows *polyhedron std-cube*

by (*simp add: std-cube-def polytope-convex-hull polytope-imp-polyhedron*)

6.2 Facets, edges, and vertices

lemma *cube-facets*:

fixes *f :: (real³) set*

shows *f face-of std-cube ∧ aff-dim f = 2 ⟷*

f = convex hull {vector[- 1, - 1, - 1], vector[- 1, - 1, 1], vector[- 1, 1, - 1], vector[- 1, 1, 1]} ∨

f = convex hull {vector[- 1, - 1, - 1], vector[- 1, - 1, 1], vector[1, - 1, - 1], vector[1, - 1, 1]} ∨

f = convex hull {vector[- 1, - 1, - 1], vector[- 1, 1, - 1], vector[1, - 1, - 1], vector[1, 1, - 1]} ∨

f = convex hull {vector[- 1, - 1, 1], vector[- 1, 1, 1], vector[1, - 1, 1], vector[1, 1, 1]} ∨

f = convex hull {vector[- 1, 1, - 1], vector[- 1, 1, 1], vector[1, 1, - 1], vector[1, 1, 1]} ∨

f = convex hull {vector[1, - 1, - 1], vector[1, - 1, 1], vector[1, 1, - 1], vector[1, 1, 1]}

unfolding *std-cube-def*

apply (*simp only: compute-faces-2 finite-insert finite.emptyI*)

apply (*simp only: compute-faces-2-step-1 bex-empty simp-thms(31)*)

apply (*simp only: compute-faces-2-step-2 bex-empty simp-thms(31)*)

apply (*simp only: empty-iff ball-insert bex-empty bex-simps(5) ball-simps(5) simp-thms(21,31)*)

apply (*simp add: vector3-sub vector3-cross vector3-eq-0 vector3-dot Let-def*)

apply (*simp only: insert-commute*)

by *linarith*

lemma *cube-facet1-edges*:

defines *v1 ≡ vector[-1, -1, -1] :: real³*

and *v2 ≡ vector[-1, -1, 1] :: real³*

and *v3 ≡ vector[-1, 1, -1] :: real³*

and *v4 ≡ vector[-1, 1, 1] :: real³*

and *n ≡ vector[1, 0, 0] :: real³*

shows (*e face-of (convex hull {v1, v2, v3, v4}) ∧ aff-dim e = 1*) =

(e = convex hull {v1, v2} ∨

```

    e = convex hull {v1, v3} ∨
    e = convex hull {v2, v4} ∨
    e = convex hull {v3, v4})
proof -
  have n-nz: n ≠ 0
    unfolding n-def by (simp add: vector3-eq-0)
  have plane:  $\bigwedge x. x \in \{v1, v2, v3, v4\} \implies n \cdot x = -1$ 
    unfolding assms using vector3-dot by auto
  show ?thesis
    by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemma cube-facet2-edges:
  defines v1  $\equiv$  vector[-1, -1, -1] :: real^3
    and v2  $\equiv$  vector[-1, -1, 1] :: real^3
    and v3  $\equiv$  vector[1, -1, -1] :: real^3
    and v4  $\equiv$  vector[1, -1, 1] :: real^3
    and n  $\equiv$  vector[0, 1, 0] :: real^3
  shows (e face-of (convex hull {v1, v2, v3, v4})  $\wedge$  aff-dim e = 1) =
    (e = convex hull {v1, v2} ∨
    e = convex hull {v1, v3} ∨
    e = convex hull {v2, v4} ∨
    e = convex hull {v3, v4})
proof -
  have n-nz: n ≠ 0
    unfolding n-def by (simp add: vector3-eq-0)
  have plane:  $\bigwedge x. x \in \{v1, v2, v3, v4\} \implies n \cdot x = -1$ 
    unfolding assms using vector3-dot by auto
  show ?thesis
    by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemma cube-facet3-edges:
  defines v1  $\equiv$  vector[-1, -1, -1] :: real^3
    and v2  $\equiv$  vector[-1, 1, -1] :: real^3
    and v3  $\equiv$  vector[1, -1, -1] :: real^3
    and v4  $\equiv$  vector[1, 1, -1] :: real^3
    and n  $\equiv$  vector[0, 0, 1] :: real^3
  shows (e face-of (convex hull {v1, v2, v3, v4})  $\wedge$  aff-dim e = 1) =
    (e = convex hull {v1, v2} ∨
    e = convex hull {v1, v3} ∨
    e = convex hull {v2, v4} ∨
    e = convex hull {v3, v4})
proof -
  have n-nz: n ≠ 0
    unfolding n-def by (simp add: vector3-eq-0)
  have plane:  $\bigwedge x. x \in \{v1, v2, v3, v4\} \implies n \cdot x = -1$ 
    unfolding assms using vector3-dot by auto
  show ?thesis

```

by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemma cube-facet4-edges:

defines $v1 \equiv \text{vector}[-1, -1, 1] :: \text{real}^3$
and $v2 \equiv \text{vector}[-1, 1, 1] :: \text{real}^3$
and $v3 \equiv \text{vector}[1, -1, 1] :: \text{real}^3$
and $v4 \equiv \text{vector}[1, 1, 1] :: \text{real}^3$
and $n \equiv \text{vector}[0, 0, 1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee$
 $e = \text{convex hull } \{v1, v3\} \vee$
 $e = \text{convex hull } \{v2, v4\} \vee$
 $e = \text{convex hull } \{v3, v4\})$

proof –

have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** (simp add: vector3-eq-0)
have plane: $\bigwedge x. x \in \{v1, v2, v3, v4\} \implies n \cdot x = 1$
unfolding assms **using** vector3-dot **by** auto
show ?thesis
by (facet-edges plane: plane n-nz: n-nz defs: assms)

qed

lemma cube-facet5-edges:

defines $v1 \equiv \text{vector}[-1, 1, -1] :: \text{real}^3$
and $v2 \equiv \text{vector}[-1, 1, 1] :: \text{real}^3$
and $v3 \equiv \text{vector}[1, 1, -1] :: \text{real}^3$
and $v4 \equiv \text{vector}[1, 1, 1] :: \text{real}^3$
and $n \equiv \text{vector}[0, 1, 0] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee$
 $e = \text{convex hull } \{v1, v3\} \vee$
 $e = \text{convex hull } \{v2, v4\} \vee$
 $e = \text{convex hull } \{v3, v4\})$

proof –

have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** (simp add: vector3-eq-0)
have plane: $\bigwedge x. x \in \{v1, v2, v3, v4\} \implies n \cdot x = 1$
unfolding assms **using** vector3-dot **by** auto
show ?thesis
by (facet-edges plane: plane n-nz: n-nz defs: assms)

qed

lemma cube-facet6-edges:

defines $v1 \equiv \text{vector}[1, -1, -1] :: \text{real}^3$
and $v2 \equiv \text{vector}[1, -1, 1] :: \text{real}^3$
and $v3 \equiv \text{vector}[1, 1, -1] :: \text{real}^3$
and $v4 \equiv \text{vector}[1, 1, 1] :: \text{real}^3$
and $n \equiv \text{vector}[1, 0, 0] :: \text{real}^3$

shows (e face-of ($\text{convex hull } \{v1, v2, v3, v4\}$) $\wedge$ $\text{aff-dim } e = 1$) =
 $(e = \text{convex hull } \{v1, v2\} \vee$
 $e = \text{convex hull } \{v1, v3\} \vee$
 $e = \text{convex hull } \{v2, v4\} \vee$
 $e = \text{convex hull } \{v3, v4\})$

proof –

have $n\text{-nz}$: $n \neq 0$
unfolding $n\text{-def}$ **by** ($\text{simp add: vector3-eq-0}$)
have plane : $\bigwedge x. x \in \{v1, v2, v3, v4\} \implies n \cdot x = 1$
unfolding assms **using** vector3-dot **by** auto
show $?thesis$
by ($\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms}$)

qed

lemmas $\text{cube-facet-edges} =$
 $\text{cube-facet1-edges cube-facet2-edges cube-facet3-edges}$
 $\text{cube-facet4-edges cube-facet5-edges cube-facet6-edges}$

lemma $\text{cube-edges-per-face}$:

assumes f face-of std-cube

and $\text{aff-dim } f = 2$

shows $\text{card } \{e. e \text{ face-of } \text{std-cube} \wedge \text{aff-dim } e = 1 \wedge e \subseteq f\} = 4$

using $\text{conjI}[OF \text{ assms}]$ **unfolding** $\text{cube-facets face-edge-set-eq}[OF \text{ assms}(1)]$

apply ($\text{elim disjE forw-subst}$)

unfolding $\text{cube-facet-edges set-cases card-4-iff2}$

$\text{segment-convex-hull[symmetric] closed-segment-eq}$

by ($\text{simp-all add: vector3-eq doubleton-eq-iff}$)

lemma cube-edges :

e face-of $\text{std-cube} \wedge \text{aff-dim } e = 1 \longleftrightarrow$

$e = \text{convex hull } \{\text{vector}[-1, -1, -1], \text{vector}[-1, -1, 1]\} \vee$

$e = \text{convex hull } \{\text{vector}[-1, -1, -1], \text{vector}[-1, 1, -1]\} \vee$

$e = \text{convex hull } \{\text{vector}[-1, -1, -1], \text{vector}[1, -1, -1]\} \vee$

$e = \text{convex hull } \{\text{vector}[-1, -1, 1], \text{vector}[-1, 1, 1]\} \vee$

$e = \text{convex hull } \{\text{vector}[-1, -1, 1], \text{vector}[1, -1, 1]\} \vee$

$e = \text{convex hull } \{\text{vector}[-1, 1, -1], \text{vector}[-1, 1, 1]\} \vee$

$e = \text{convex hull } \{\text{vector}[-1, 1, -1], \text{vector}[1, 1, -1]\} \vee$

$e = \text{convex hull } \{\text{vector}[-1, 1, 1], \text{vector}[1, 1, 1]\} \vee$

$e = \text{convex hull } \{\text{vector}[1, -1, -1], \text{vector}[1, -1, 1]\} \vee$

$e = \text{convex hull } \{\text{vector}[1, -1, -1], \text{vector}[1, 1, -1]\} \vee$

$e = \text{convex hull } \{\text{vector}[1, -1, 1], \text{vector}[1, 1, 1]\} \vee$

$e = \text{convex hull } \{\text{vector}[1, 1, -1], \text{vector}[1, 1, 1]\}$

(**is** $?lhs \longleftrightarrow ?rhs$)

proof

assume $*$: $?lhs$

then obtain f **where** $f\text{-facts}$: f face-of $\text{std-cube} \wedge \text{aff-dim } f = 2$ e face-of $f \wedge$
 $\text{aff-dim } e = 1$

using $\text{cube-polyhedron cube-fulldim edge-belongs-to-face}$

by blast

```

show ?rhs
  using f-facts unfolding cube-facets
  apply (elim disjE) apply (all ⟨hypsubst-thin⟩)
  unfolding cube-facet-edges
  by linarith+
next
  assume ?rhs
  then show ?lhs
    using cube-facets cube-facet-edges
    by (smt (verit, ccfv-threshold) face-of-trans insertI1 insert-commute mem-Collect-eq)
qed

```

lemma *cube-vertices*:

```

shows (v face-of std-cube ∧ aff-dim v = 0) =
  (v = {vector [1, 1, 1]} ∨
   v = {vector [1, 1, -1]} ∨
   v = {vector [1, -1, 1]} ∨
   v = {vector [1, -1, -1]} ∨
   v = {vector [-1, 1, 1]} ∨
   v = {vector [-1, 1, -1]} ∨
   v = {vector [-1, -1, 1]} ∨
   v = {vector [-1, -1, -1]})
(is ?lhs = ?rhs)
proof
  assume *: ?lhs
  obtain e where e face-of std-cube aff-dim e = 1 v face-of e
    using * cube-polyhedron cube-fulldim vertex-belongs-to-edge
    by blast
  show ?rhs
    using conjI[OF ⟨e face-of std-cube⟩ ⟨aff-dim e = 1⟩]
    unfolding cube-edges
    apply (elim disjE)
    using vertices-of-edge[OF - ⟨v face-of e⟩] *
    by metis+
next
  assume ?rhs
  then show ?lhs
    apply (elim disjE)
    by (meson aff-dim-sing extreme-point-of-convex-hull-2 face-of-singleton
        face-of-trans cube-edges)+
qed

```

lemma *cube-edges-per-vertex*:

```

assumes v face-of std-cube
  and aff-dim v = 0
shows card {e. e face-of std-cube ∧ aff-dim e = 1 ∧ v ⊆ e} = 3
unfolding vertex-edge-set-eq[OF assms(1)]
using conjI[OF assms]
unfolding cube-vertices conj-assoc[symmetric] cube-edges

```

by (elim disjE forw-subst) edges-per-vertex+

6.3 Regularity

lemma cube-congruent-edges:

assumes $e1$ face-of std-cube $\wedge$ aff-dim $e1 = 1$

and $e2$ face-of std-cube $\wedge$ aff-dim $e2 = 1$

shows $e1$ congruent $e2$

proof –

let ?edges =

{convex hull {(vector[- 1, - 1, - 1]::real^3), vector[- 1, - 1, 1]},
convex hull {vector[- 1, - 1, - 1], vector[- 1, 1, - 1]},
convex hull {vector[- 1, - 1, - 1], vector[1, - 1, - 1]},
convex hull {vector[- 1, - 1, 1], vector[- 1, 1, 1]},
convex hull {vector[- 1, - 1, 1], vector[1, - 1, 1]},
convex hull {vector[- 1, 1, - 1], vector[- 1, 1, 1]},
convex hull {vector[- 1, 1, - 1], vector[1, 1, - 1]},
convex hull {vector[- 1, 1, 1], vector[1, 1, 1]},
convex hull {vector[1, - 1, - 1], vector[1, - 1, 1]},
convex hull {vector[1, - 1, - 1], vector[1, 1, - 1]},
convex hull {vector[1, - 1, 1], vector[1, 1, 1]},
convex hull {vector[1, 1, - 1], vector[1, 1, 1]}}

have $e1 \in ?edges$ $e2 \in ?edges$

using assms **unfolding** cube-edges

by auto

moreover **have** $\forall e \in ?edges. \text{convex hull } \{ \text{vector}[- 1, - 1, - 1], \text{vector}[- 1, - 1, 1] \}$ congruent e

apply (simp only: ball-simps(7) Set.ball-empty simp-thms(21))

unfolding conj-assoc[symmetric] segment-convex-hull[symmetric]

apply (rule conjI)+

apply (all (rule congruent-segments))

unfolding dist-norm vector3-sub

apply (simp-all add: norm-eq)

unfolding vector3-dot

by linarith+

ultimately show ?thesis

using congruent-set **by** meson

qed

lemma cube-congruent-faces:

assumes $f1$ face-of std-cube $\wedge$ aff-dim $f1 = 2$

and $f2$ face-of std-cube $\wedge$ aff-dim $f2 = 2$

shows $f1$ congruent $f2$

proof –

let ?faces =

{convex hull {(vector[- 1, - 1, - 1]::real^3), vector[- 1, - 1, 1], vector[- 1, 1, - 1], vector[- 1, 1, 1]},
convex hull {vector[- 1, - 1, - 1], vector[- 1, - 1, 1], vector[1, - 1, - 1], vector[1, - 1, 1]},
convex hull {vector[- 1, 1, - 1], vector[- 1, 1, 1], vector[1, 1, - 1], vector[1, 1, 1]},
convex hull {vector[1, - 1, - 1], vector[1, - 1, 1], vector[1, 1, - 1], vector[1, 1, 1]}}

```

      convex hull {vector[- 1, - 1, - 1], vector[- 1, 1, - 1], vector[1, - 1, -
1], vector[1, 1, - 1]},
      convex hull {vector[- 1, - 1, 1], vector[- 1, 1, 1], vector[1, - 1, 1], vector[1,
1, 1]},
      convex hull {vector[- 1, 1, - 1], vector[- 1, 1, 1], vector[1, 1, - 1], vector[1,
1, 1]},
      convex hull {vector[1, - 1, - 1], vector[1, - 1, 1], vector[1, 1, - 1], vector[1,
1, 1]}
    have f1 ∈ ?faces f2 ∈ ?faces
    using assms cube-facets
    by auto
  moreover have ∀ f ∈ ?faces.
    convex hull {vector[- 1, - 1, - 1], vector[- 1, - 1, 1], vector[- 1, 1,
- 1], vector[- 1, 1, 1]} congruent f
    apply (simp only: ball-simps(7) Set.ball-empty simp-thms(21))
    unfolding conj-assoc[symmetric]
    by (rule conjI)+ show-congruence+
  ultimately show ?thesis
    using congruent-set by meson
qed

```

lemma *cube-facet1-equangular*:

```

defines v1 ≡ vector[-1, -1, -1] :: real^3
defines v2 ≡ vector[-1, -1, 1] :: real^3
defines v3 ≡ vector[-1, 1, -1] :: real^3
defines v4 ≡ vector[-1, 1, 1] :: real^3
shows equiangular (convex hull {v1, v2, v3, v4}) (pi / 2)
unfolding equiangular-def
proof (clarify)
  fix e1 e2 v
  assume *: e1 face-of convex hull {v1, v2, v3, v4} aff-dim e1 = 1
            e2 face-of convex hull {v1, v2, v3, v4} aff-dim e2 = 1
            e1 ≠ e2 v extreme-point-of e1 v extreme-point-of e2
  have (v = v1 ∧ ((e1 = convex hull {v1, v2} ∧ e2 = convex hull {v1, v3}) ∨
    (e1 = convex hull {v1, v3} ∧ e2 = convex hull {v1, v2}))) ∨
    (v = v2 ∧ ((e1 = convex hull {v2, v1} ∧ e2 = convex hull {v2, v4}) ∨
    (e1 = convex hull {v2, v4} ∧ e2 = convex hull {v2, v1}))) ∨
    (v = v3 ∧ ((e1 = convex hull {v3, v1} ∧ e2 = convex hull {v3, v4}) ∨
    (e1 = convex hull {v3, v4} ∧ e2 = convex hull {v3, v1}))) ∨
    (v = v4 ∧ ((e1 = convex hull {v4, v2} ∧ e2 = convex hull {v4, v3}) ∨
    (e1 = convex hull {v4, v3} ∧ e2 = convex hull {v4, v2})))
  using conjI[OF *(1,2)] conjI[OF *(3,4)] *(5-7) unfolding assms cube-facet1-edges
  by (smt (verit, best) extreme-point-of-convex-hull-2 insert-commute vector3-eq)
  moreover have vangle (v2 - v1) (v3 - v1) = pi / 2
    vangle (v1 - v2) (v4 - v2) = pi / 2
    vangle (v1 - v3) (v4 - v3) = pi / 2
    vangle (v2 - v4) (v3 - v4) = pi / 2
  unfolding assms vector3-sub norm-eq-sqrt-inner vector3-dot vangle-def
  by (simp-all add: vector3-eq-0)

```

```

ultimately show ( $\exists a\ b. e1 = \text{convex hull } \{v, a\} \wedge e2 = \text{convex hull } \{v, b\} \wedge$ 
                   $\text{vangle } (a - v) (b - v) = \text{pi} / 2$ )
  using vangle-commute
  by metis
qed

lemma cube-equiangular:
  assumes  $f \text{ face-of std-cube} \wedge \text{aff-dim } f = 2$ 
  shows equiangular  $f$  ( $\text{pi} / 2$ )
proof -
  define f1 where  $f1 = \text{convex hull } \{\text{vector}[-1, -1, -1]::\text{real}^3, \text{vector}[-1,$ 
 $-1, 1], \text{vector}[-1, 1, -1], \text{vector}[-1, 1, 1]\}$ 
  have  $f1 \text{ face-of std-cube} \wedge \text{aff-dim } f1 = 2$ 
  using f1-def cube-facets by blast
  then have  $f1 \text{ congruent } f$ 
  using assms cube-congruent-faces by blast
  then show ?thesis
  using equiangular-congruent cube-facet1-equiangular f1-def
  by blast
qed

end

```

7 Octahedron

```

theory Octahedron
  imports Computation
begin

```

7.1 Definition

```

definition std-octahedron :: (real, 3) vec set
  where  $\text{std-octahedron} \equiv \text{convex hull}$ 
 $\{\text{vector } [1, 0, 0], \text{vector } [-1, 0, 0],$ 
 $\text{vector } [0, 0, 1], \text{vector } [0, 0, -1],$ 
 $\text{vector } [0, 1, 0], \text{vector } [0, -1, 0]\}$ 

lemma octahedron-fulldim:
  shows  $\text{aff-dim } \text{std-octahedron} = 3$ 
  unfolding std-octahedron-def
  apply (rule polytope-3D)
  by (simp add: vector3-sub vector3-cross vector3-dot vector3-eq-0)+

lemma octahedron-polyhedron:
  shows polyhedron std-octahedron
  by (simp add: std-octahedron-def polytope-convex-hull polytope-imp-polyhedron)

```

7.2 Facets, edges, and vertices

lemma *octahedron-facets:*

fixes $f :: (\text{real}^3) \text{ set}$
shows $f \text{ face-of std-octahedron} \wedge \text{aff-dim } f = 2 \iff$
 $f = \text{convex hull } \{\text{vector}[-1, 0, 0], \text{vector}[0, -1, 0], \text{vector}[0, 0, -1]\} \vee$
 $f = \text{convex hull } \{\text{vector}[-1, 0, 0], \text{vector}[0, -1, 0], \text{vector}[0, 0, 1]\} \vee$
 $f = \text{convex hull } \{\text{vector}[-1, 0, 0], \text{vector}[0, 1, 0], \text{vector}[0, 0, -1]\} \vee$
 $f = \text{convex hull } \{\text{vector}[-1, 0, 0], \text{vector}[0, 1, 0], \text{vector}[0, 0, 1]\} \vee$
 $f = \text{convex hull } \{\text{vector}[1, 0, 0], \text{vector}[0, -1, 0], \text{vector}[0, 0, -1]\} \vee$
 $f = \text{convex hull } \{\text{vector}[1, 0, 0], \text{vector}[0, -1, 0], \text{vector}[0, 0, 1]\} \vee$
 $f = \text{convex hull } \{\text{vector}[1, 0, 0], \text{vector}[0, 1, 0], \text{vector}[0, 0, -1]\} \vee$
 $f = \text{convex hull } \{\text{vector}[1, 0, 0], \text{vector}[0, 1, 0], \text{vector}[0, 0, 1]\}$
unfolding *std-octahedron-def*
apply (*simp only: compute-faces-2 finite-insert finite.emptyI*)
apply (*simp only: compute-faces-2-step-1 bex-empty simp-thms(31)*)
apply (*simp only: compute-faces-2-step-2 bex-empty simp-thms(31)*)
apply (*simp only: empty-iff ball-insert bex-empty bex-simps(5) ball-simps(5)*
simp-thms(21,31))
apply (*simp add: vector3-sub vector3-cross vector3-eq-0 vector3-dot Let-def*)
apply (*simp only: insert-commute*)
by *linarith*

lemma *octahedron-facet1-edges:*

defines $v1 \equiv \text{vector}[-1, 0, 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[0, -1, 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[0, 0, -1] :: \text{real}^3$
and $n \equiv \text{vector}[1, 1, 1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding *n-def* **by** (*simp add: vector3-eq-0*)
have *plane*: $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = -1$
by (*auto simp add: assms vector3-dot*)
have *finite*: *finite* $\{v1, v2, v3\}$
using *assms* **by** *blast*
show *?thesis*
by (*facet-edges plane: plane n-nz: n-nz defs: assms*)
qed

lemma *octahedron-facet2-edges:*

defines $v1 \equiv \text{vector}[-1, 0, 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[0, -1, 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[0, 0, 1] :: \text{real}^3$
and $n \equiv \text{vector}[-1, -1, 1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$

```

proof –
  have  $n\text{-nz}: n \neq 0$ 
    unfolding  $n\text{-def}$  by (simp add: vector3-eq-0)
  have  $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = 1$ 
    by (auto simp add: assms vector3-dot)
  have  $\text{finite}: \text{finite } \{v1, v2, v3\}$ 
    using assms by blast
  show ?thesis
    by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemma octahedron-facet3-edges:
  defines  $v1 \equiv \text{vector}[-1, 0, 0] :: \text{real}^3$ 
    and  $v2 \equiv \text{vector}[0, 1, 0] :: \text{real}^3$ 
    and  $v3 \equiv \text{vector}[0, 0, -1] :: \text{real}^3$ 
    and  $n \equiv \text{vector}[-1, 1, -1] :: \text{real}^3$ 
  shows  $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$ 
     $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$ 
proof –
  have  $n\text{-nz}: n \neq 0$ 
    unfolding  $n\text{-def}$  by (simp add: vector3-eq-0)
  have  $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = 1$ 
    by (auto simp add: assms vector3-dot)
  have  $\text{finite}: \text{finite } \{v1, v2, v3\}$ 
    using assms by blast
  show ?thesis
    by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemma octahedron-facet4-edges:
  defines  $v1 \equiv \text{vector}[-1, 0, 0] :: \text{real}^3$ 
    and  $v2 \equiv \text{vector}[0, 1, 0] :: \text{real}^3$ 
    and  $v3 \equiv \text{vector}[0, 0, 1] :: \text{real}^3$ 
    and  $n \equiv \text{vector}[1, -1, -1] :: \text{real}^3$ 
  shows  $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$ 
     $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$ 
proof –
  have  $n\text{-nz}: n \neq 0$ 
    unfolding  $n\text{-def}$  by (simp add: vector3-eq-0)
  have  $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = -1$ 
    by (auto simp add: assms vector3-dot)
  have  $\text{finite}: \text{finite } \{v1, v2, v3\}$ 
    using assms by blast
  show ?thesis
    by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

```

lemma *octahedron-facet5-edges*:
defines $v1 \equiv \text{vector}[1, 0, 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[0, -1, 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[0, 0, -1] :: \text{real}^3$
and $n \equiv \text{vector}[1, -1, -1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = 1$
by $(\text{auto simp add: assms vector3-dot})$
have $\text{finite}: \text{finite } \{v1, v2, v3\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$
qed

lemma *octahedron-facet6-edges*:
defines $v1 \equiv \text{vector}[1, 0, 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[0, -1, 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[0, 0, 1] :: \text{real}^3$
and $n \equiv \text{vector}[-1, 1, -1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = -1$
by $(\text{auto simp add: assms vector3-dot})$
have $\text{finite}: \text{finite } \{v1, v2, v3\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$
qed

lemma *octahedron-facet7-edges*:
defines $v1 \equiv \text{vector}[1, 0, 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[0, 1, 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[0, 0, -1] :: \text{real}^3$
and $n \equiv \text{vector}[-1, -1, 1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0})$

```

have plane:  $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = -1$ 
  by (auto simp add: assms vector3-dot)
have finite: finite {v1, v2, v3}
  using assms by blast
show ?thesis
  by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

```

```

lemma octahedron-facet8-edges:
  defines v1  $\equiv$  vector[1, 0, 0] :: real^3
    and v2  $\equiv$  vector[0, 1, 0] :: real^3
    and v3  $\equiv$  vector[0, 0, 1] :: real^3
    and n  $\equiv$  vector[1, 1, 1] :: real^3
  shows (e face-of (convex hull {v1, v2, v3})  $\wedge$  aff-dim e = 1) =
    (e = convex hull {v1, v2}  $\vee$  e = convex hull {v1, v3}  $\vee$  e = convex hull
    {v2, v3})
proof -
  have n-nz: n  $\neq$  0
    unfolding n-def by (simp add: vector3-eq-0)
  have plane:  $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = 1$ 
    by (auto simp add: assms vector3-dot)
  have finite: finite {v1, v2, v3}
    using assms by blast
  show ?thesis
    by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

```

```

lemmas octahedron-facet-edges =
  octahedron-facet1-edges octahedron-facet2-edges
  octahedron-facet3-edges octahedron-facet4-edges octahedron-facet5-edges
  octahedron-facet6-edges octahedron-facet7-edges octahedron-facet8-edges

```

```

lemma octahedron-edges-per-face:
  assumes f face-of std-octahedron
    and aff-dim f = 2
  shows card {e. e face-of std-octahedron  $\wedge$  aff-dim e = 1  $\wedge$  e  $\subseteq$  f} = 3
  using conjI[OF assms] unfolding octahedron-facets face-edge-set-eq[OF assms(1)]
  apply (elim disjE forw-subst)
  unfolding octahedron-facet-edges set-cases card-3-iff2
    segment-convex-hull[symmetric] closed-segment-eq
  by (simp-all add: vector3-eq doubleton-eq-iff)

```

```

lemma octahedron-edges:
  e face-of std-octahedron  $\wedge$  aff-dim e = 1  $\longleftrightarrow$ 
  e = convex hull {vector[- 1, 0, 0], vector[0, - 1, 0]}  $\vee$ 
  e = convex hull {vector[- 1, 0, 0], vector[0, 1, 0]}  $\vee$ 
  e = convex hull {vector[- 1, 0, 0], vector[0, 0, - 1]}  $\vee$ 
  e = convex hull {vector[- 1, 0, 0], vector[0, 0, 1]}  $\vee$ 
  e = convex hull {vector[1, 0, 0], vector[0, - 1, 0]}  $\vee$ 

```

```

    e = convex hull {vector[1, 0, 0], vector[0, 1, 0]} ∨
    e = convex hull {vector[1, 0, 0], vector[0, 0, -1]} ∨
    e = convex hull {vector[1, 0, 0], vector[0, 0, 1]} ∨
    e = convex hull {vector[0, -1, 0], vector[0, 0, -1]} ∨
    e = convex hull {vector[0, -1, 0], vector[0, 0, 1]} ∨
    e = convex hull {vector[0, 1, 0], vector[0, 0, -1]} ∨
    e = convex hull {vector[0, 1, 0], vector[0, 0, 1]}
    (is ?lhs ⟷ ?rhs)
  proof
    assume ?lhs
    then obtain f where f-facts: f face-of std-octahedron ∧ aff-dim f = 2 e face-of
    f ∧ aff-dim e = 1
      using octahedron-polyhedron octahedron-fulldim edge-belongs-to-face
      by blast
    show ?rhs
      using f-facts unfolding octahedron-facets
      apply (elim disjE) apply (all ‹hypsubst-thin›)
      unfolding octahedron-facet-edges
      by linarith+
  next
    assume ?rhs
    then show ?lhs
      using octahedron-facets octahedron-facet-edges
      by (smt (verit, ccfv-threshold) face-of-trans insertI1 insert-commute mem-Collect-eq)
  qed

lemma octahedron-vertices:
  shows (v face-of std-octahedron ∧ aff-dim v = 0) =
    (v = {vector [1, 0, 0]} ∨
     v = {vector [-1, 0, 0]} ∨
     v = {vector [0, 0, 1]} ∨
     v = {vector [0, 0, -1]} ∨
     v = {vector [0, 1, 0]} ∨
     v = {vector [0, -1, 0]})
  (is ?lhs = ?rhs)
proof
  assume *: ?lhs
  then obtain e where e-facts: e face-of std-octahedron ∧ aff-dim e = 1 v face-of
  e
    using octahedron-polyhedron octahedron-fulldim vertex-belongs-to-edge
    by blast
  show ?rhs
    using e-facts(1)
    unfolding octahedron-edges
    apply (elim disjE)
    using vertices-of-edge[OF - ‹v face-of e›] *
    by metis+
  next
    assume ?rhs

```

```

then show ?lhs
  apply (elim disjE)
  by (meson aff-dim-sing extreme-point-of-convex-hull-2 face-of-singleton
      face-of-trans octahedron-edges)+
qed

```

```

lemma octahedron-edges-per-vertex:
  assumes v face-of std-octahedron
    and aff-dim v = 0
  shows card {e. e face-of std-octahedron ∧ aff-dim e = 1 ∧ v ⊆ e} = 4
  unfolding vertex-edge-set-eq[OF assms(1)]
  using conjI[OF assms]
  unfolding octahedron-vertices conj-assoc[symmetric] octahedron-edges
  by (elim disjE forw-subst) edges-per-vertex+

```

7.3 Regularity

```

lemma octahedron-congruent-edges:
  assumes e1 face-of std-octahedron ∧ aff-dim e1 = 1
    and e2 face-of std-octahedron ∧ aff-dim e2 = 1
  shows e1 congruent e2
proof -
  let ?edges =
    {convex hull {(vector[- 1, 0, 0]::real^3), vector[0, - 1, 0]},
     convex hull {vector[- 1, 0, 0], vector[0, 1, 0]},
     convex hull {vector[- 1, 0, 0], vector[0, 0, - 1]},
     convex hull {vector[- 1, 0, 0], vector[0, 0, 1]},
     convex hull {vector[1, 0, 0], vector[0, - 1, 0]},
     convex hull {vector[1, 0, 0], vector[0, 1, 0]},
     convex hull {vector[1, 0, 0], vector[0, 0, - 1]},
     convex hull {vector[1, 0, 0], vector[0, 0, 1]},
     convex hull {vector[0, - 1, 0], vector[0, 0, - 1]},
     convex hull {vector[0, - 1, 0], vector[0, 0, 1]},
     convex hull {vector[0, 1, 0], vector[0, 0, - 1]},
     convex hull {vector[0, 1, 0], vector[0, 0, 1]}}
  have e1 ∈ ?edges e2 ∈ ?edges
  using assms unfolding octahedron-edges
  by auto
  moreover have ∀ e ∈ ?edges. convex hull {vector[- 1, 0, 0], vector[0, - 1, 0]}
    congruent e
  apply (simp only: ball-simps(7) Set.ball-empty simp-thms(21))
  unfolding conj-assoc[symmetric] segment-convex-hull[symmetric]
  apply (rule conjI)+
  apply (all (rule congruent-segments))
  unfolding dist-norm vector3-sub
  apply (simp-all add: norm-eq)
  unfolding vector3-dot
  by linarith+
ultimately show ?thesis

```

using congruent-set by meson
qed

lemma octahedron-congruent-faces:

assumes $f1$ face-of std-octahedron $\wedge$ aff-dim $f1 = 2$
and $f2$ face-of std-octahedron $\wedge$ aff-dim $f2 = 2$
shows $f1$ congruent $f2$

proof –

let ?faces =

{convex hull {(vector[- 1, 0, 0]::real^3), vector[0, - 1, 0], vector[0, 0, - 1]},
convex hull {vector[- 1, 0, 0], vector[0, - 1, 0], vector[0, 0, 1]},
convex hull {vector[- 1, 0, 0], vector[0, 1, 0], vector[0, 0, - 1]},
convex hull {vector[- 1, 0, 0], vector[0, 1, 0], vector[0, 0, 1]},
convex hull {vector[1, 0, 0], vector[0, - 1, 0], vector[0, 0, - 1]},
convex hull {vector[1, 0, 0], vector[0, - 1, 0], vector[0, 0, 1]},
convex hull {vector[1, 0, 0], vector[0, 1, 0], vector[0, 0, - 1]},
convex hull {vector[1, 0, 0], vector[0, 1, 0], vector[0, 0, 1]}}

have $f1 \in ?faces$ $f2 \in ?faces$

using assms octahedron-facets

by auto

moreover have $\forall f \in ?faces.$

convex hull {vector[- 1, 0, 0], vector[0, - 1, 0], vector[0, 0, - 1]}

congruent f

apply (simp only: ball-simps(7) Set.ball-empty simp-thms(21))

unfolding conj-assoc[symmetric]

by (rule conjI)+ show-congruence+

ultimately show ?thesis

using congruent-set by meson

qed

lemma octahedron-facet1-equiangular:

defines $v1 \equiv \text{vector}[-1, 0, 0] :: \text{real}^3$

defines $v2 \equiv \text{vector}[0, -1, 0] :: \text{real}^3$

defines $v3 \equiv \text{vector}[0, 0, -1] :: \text{real}^3$

shows equiangular (convex hull { $v1, v2, v3$ }) ($\pi / 3$)

unfolding equiangular-def

proof (clarify)

fix $e1$ $e2$ v

assume *: $e1$ face-of convex hull { $v1, v2, v3$ } aff-dim $e1 = 1$

$e2$ face-of convex hull { $v1, v2, v3$ } aff-dim $e2 = 1$

$e1 \neq e2$ v extreme-point-of $e1$ v extreme-point-of $e2$

have ($v = v1 \wedge ((e1 = \text{convex hull } \{v1, v2\} \wedge e2 = \text{convex hull } \{v1, v3\}) \vee$
($e1 = \text{convex hull } \{v1, v3\} \wedge e2 = \text{convex hull } \{v1, v2\}$))) $\vee$
($v = v2 \wedge ((e1 = \text{convex hull } \{v2, v1\} \wedge e2 = \text{convex hull } \{v2, v3\}) \vee$
($e1 = \text{convex hull } \{v2, v3\} \wedge e2 = \text{convex hull } \{v2, v1\}$))) $\vee$
($v = v3 \wedge ((e1 = \text{convex hull } \{v3, v1\} \wedge e2 = \text{convex hull } \{v3, v2\}) \vee$
($e1 = \text{convex hull } \{v3, v2\} \wedge e2 = \text{convex hull } \{v3, v1\}$)))

using conjI[OF *(1,2)] conjI[OF *(3,4)] *(5-7) unfolding assms octahedron-facet1-edges

```

    by (smt (verit, best) extreme-point-of-convex-hull-2 insert-commute vector3-eq)
  moreover have  $\text{vangle } (v2 - v1) (v3 - v1) = \pi / 3$ 
    vangle  $(v1 - v2) (v3 - v2) = \pi / 3$ 
    vangle  $(v1 - v3) (v2 - v3) = \pi / 3$ 
  unfolding assms vector3-sub norm-eq-sqrt-inner vector3-dot vangle-def
  by (simp-all add: vector3-eq-0)
  ultimately show  $(\exists a b. e1 = \text{convex hull } \{v, a\} \wedge e2 = \text{convex hull } \{v, b\} \wedge$ 
     $\text{vangle } (a - v) (b - v) = \pi / 3)$ 
    using vangle-commute
    by metis
qed

```

lemma *octahedron-equiangular*:

```

  assumes  $f \text{ face-of std-octahedron} \wedge \text{aff-dim } f = 2$ 
  shows equiangular  $f (\pi / 3)$ 
proof -
  define f1 where  $f1 = \text{convex hull } \{\text{vector}[-1, 0, 0]::\text{real}^3, \text{vector}[0, -1, 0],$ 
     $\text{vector}[0, 0, -1]\}$ 
  have  $f1 \text{ face-of std-octahedron} \wedge \text{aff-dim } f1 = 2$ 
    using f1-def octahedron-facets by blast
  then have  $f1 \text{ congruent } f$ 
    using assms octahedron-congruent-faces by blast
  then show ?thesis
    using equiangular-congruent octahedron-facet1-equiangular f1-def
    by blast
qed
end

```

8 Dodecahedron

theory *Dodecahedron*

```

  imports
    Computation
    HOL-Library.Quadratic-Discriminant
begin

```

8.1 Definition

definition *std-dodecahedron* :: $(\text{real}, 3) \text{ vec set}$

```

  where  $\text{std-dodecahedron} \equiv \text{convex hull}$ 
     $\{\text{vector } [1, 1, 1], \text{vector } [1, 1, -1],$ 
       $\text{vector } [1, -1, 1], \text{vector } [1, -1, -1],$ 
       $\text{vector } [-1, 1, 1], \text{vector } [-1, 1, -1],$ 
       $\text{vector } [-1, -1, 1], \text{vector } [-1, -1, -1],$ 
       $\text{vector } [0, \text{inverse } \varphi, \varphi], \text{vector } [0, \text{inverse } \varphi, -\varphi],$ 
       $\text{vector } [0, -\text{inverse } \varphi, \varphi], \text{vector } [0, -\text{inverse } \varphi, -\varphi],$ 
       $\text{vector } [\text{inverse } \varphi, \varphi, 0], \text{vector } [\text{inverse } \varphi, -\varphi, 0],$ 
       $\text{vector } [-\text{inverse } \varphi, \varphi, 0], \text{vector } [-\text{inverse } \varphi, -\varphi, 0]\}$ 

```

vector $[\varphi, 0, \text{inverse } \varphi]$, vector $[-\varphi, 0, \text{inverse } \varphi]$,
vector $[\varphi, 0, -\text{inverse } \varphi]$, vector $[-\varphi, 0, -\text{inverse } \varphi]$

lemma *std-dodecahedron-explicit*:

shows *std-dodecahedron* = convex hull

{ vector[1, 1, 1],
vector[1, 1, -1],
vector[1, -1, 1],
vector[1, -1, -1],
vector[-1, 1, 1],
vector[-1, 1, -1],
vector[-1, -1, 1],
vector[-1, -1, -1],
vector[0, -1 / 2 + 1 / 2 * sqrt(5), 1 / 2 + 1 / 2 * sqrt(5)],
vector[0, -1 / 2 + 1 / 2 * sqrt(5), -1 / 2 + -1 / 2 * sqrt(5)],
vector[0, 1 / 2 + -1 / 2 * sqrt(5), 1 / 2 + 1 / 2 * sqrt(5)],
vector[0, 1 / 2 + -1 / 2 * sqrt(5), -1 / 2 + -1 / 2 * sqrt(5)],
vector[-1 / 2 + 1 / 2 * sqrt(5), 1 / 2 + 1 / 2 * sqrt(5), 0],
vector[-1 / 2 + 1 / 2 * sqrt(5), -1 / 2 + -1 / 2 * sqrt(5), 0],
vector[1 / 2 + -1 / 2 * sqrt(5), 1 / 2 + 1 / 2 * sqrt(5), 0],
vector[1 / 2 + -1 / 2 * sqrt(5), -1 / 2 + -1 / 2 * sqrt(5), 0],
vector[1 / 2 + 1 / 2 * sqrt(5), 0, -1 / 2 + 1 / 2 * sqrt(5)],
vector[-1 / 2 + -1 / 2 * sqrt(5), 0, -1 / 2 + 1 / 2 * sqrt(5)],
vector[1 / 2 + 1 / 2 * sqrt(5), 0, 1 / 2 + -1 / 2 * sqrt(5)],
vector[-1 / 2 + -1 / 2 * sqrt(5), 0, 1 / 2 + -1 / 2 * sqrt(5)] }

proof –

have *inverse* $\varphi = 2 / (1 + \text{sqrt } 5)$ **by** *simp*

also have $\dots = (2 * (1 - \text{sqrt } 5)) / ((1 + \text{sqrt } 5) * (1 - \text{sqrt } 5))$

by (*metis mult.commute nonzero-mult-divide-mult-cancel-right2 numeral-eq-one-iff*
real-sqrt-eq-1-iff right-minus-eq

verit-eq-simplify(12))

also have $\dots = (2 - 2 * \text{sqrt } 5) / (-4)$

using *square-diff-square-factored[of 1 sqrt 5]* **by** *auto*

finally have *inverse* $\varphi = -1/2 + 1/2 * \text{sqrt } 5$

by *argo*

moreover have *golden*: $\varphi = 1/2 + 1/2 * \text{sqrt } 5$

by *auto*

ultimately have *golden-inverse*: *inverse* $(1/2 + 1/2 * \text{sqrt } 5) = -1/2 + 1/2$
 $* \text{sqrt } 5$

by *auto*

show *?thesis*

unfolding *std-dodecahedron-def golden golden-inverse*

by *simp*

qed

lemma *std-dodecahedron-eq*:

shows *std-dodecahedron* = convex hull

{ vector[*rat5* 1 0, *rat5* 1 0, *rat5* 1 0],
vector[*rat5* 1 0, *rat5* 1 0, *rat5* (-1) 0],

```

vector[rat5 1 0, rat5 (-1) 0, rat5 1 0],
vector[rat5 1 0, rat5 (-1) 0, rat5 (-1) 0],
vector[rat5 (-1) 0, rat5 1 0, rat5 1 0],
vector[rat5 (-1) 0, rat5 1 0, rat5 (-1) 0],
vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0],
vector[rat5 (-1) 0, rat5 (-1) 0, rat5 (-1) 0],
vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (1/2) (1/2)],
vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2)],
vector[rat5 0 0, rat5 (1/2) (-1/2), rat5 (1/2) (1/2)],
vector[rat5 0 0, rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2)],
vector[rat5 (-1/2) (1/2), rat5 (1/2) (1/2), rat5 0 0],
vector[rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0],
vector[rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1/2) (1/2)],
vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)],
vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (1/2) (-1/2)],
vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)]}
unfolding std-dodecahedron-explicit rat5-def
by (simp add: of-rat-divide of-rat-minus)

```

lemma dodecahedron-fulldim:

shows aff-dim std-dodecahedron = 3

unfolding std-dodecahedron-def

apply (rule polytope-3D)

by (simp add: vector3-sub vector3-cross vector3-dot vector3-eq-0)+

lemma dodecahedron-polyhedron:

shows polyhedron std-dodecahedron

by (simp add: std-dodecahedron-def polytope-convex-hull polytope-imp-polyhedron)

8.2 Facets, edges, and vertices

lemma dodecahedron-facets:

fixes f :: (real³) set

shows f face-of std-dodecahedron $\wedge$ aff-dim f = 2 $\longleftrightarrow$

f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)], vector[rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 (-1) 0, rat5 (-1) 0, rat5 (-1) 0], vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0]} $\vee$

f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)], vector[rat5 (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0], vector[rat5 (-1) 0, rat5 1 0, rat5 (-1) 0], vector[rat5 (-1) 0, rat5 1 0, rat5 1 0]} $\vee$

f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0], vector[rat5 (-1) 0, rat5 1 0, rat5 1 0], vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (1/2) (1/2)], vector[rat5 0 0, rat5 (1/2) (-1/2), rat5 (1/2) (1/2)]} $\vee$

f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)],

$\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2)], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2)]\} \vee$
 $f = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2)]\} \vee$
 $f = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (1/2) \ (1/2)]\} \vee$
 $f = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2)], \text{vector}[\text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2)], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0]\} \vee$
 $f = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } (1/2) \ (-1/2), \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2)]\} \vee$
 $f = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } (1/2) \ (-1/2), \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2), \text{rat5 } (1/2) \ (1/2)]\} \vee$
 $f = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2)], \text{vector}[\text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2)], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0]\} \vee$
 $f = \text{convex hull } \{\text{vector}[\text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2)], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2), \text{rat5 } (1/2) \ (1/2)], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2)], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2)]\}$
unfolding *std-dodecahedron-eq*
apply (*simp only: compute-faces-2 finite-insert finite.emptyI*)
apply (*simp only: compute-faces-2-step-1 bex-empty simp-thms(31)*)
apply (*simp only: compute-faces-2-step-2 bex-empty simp-thms(31)*)
apply (*simp only: empty-iff ball-insert bex-empty bex-simps(5) ball-simps(5) simp-thms(21,31)*)
apply (*simp only: vector3-sub vector3-cross vector3-eq-0 rat5-add rat5-sub rat5-mul rat5-eq*)
apply (*simp add: vector3-eq-0 vector3-dot rat5-add rat5-sub rat5-mul rat5-le rat5-eq rat5-eq-0 Let-def*)
— very slow, ca 600s
apply (*simp only: insert-commute*)
by *linarith*

lemma *dodecahedron-facet1-edges*:

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) (1/2)] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) (-1/2)] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } (1/2) (-1/2), \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v4 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0] :: \text{real}^3$
and $v5 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } (-2) \ 0, \text{rat5 } 1 \ (-1), \text{rat5 } 0 \ 0] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4, v5\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v5\} \vee e = \text{convex hull } \{v2, v4\} \vee e = \text{convex hull } \{v3, v4\} \vee e = \text{convex hull } \{v3, v5\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0 rat5-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } 1 \ 1$
unfolding assms
by $(\text{auto simp add: vector3-dot rat5-mul rat5-add})$
have $\text{finite}: \text{finite } \{v1, v2, v3, v4, v5\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$
qed

lemma *dodecahedron-facet2-edges*:

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) (1/2)] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) (-1/2)] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } (1/2) (-1/2), \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v4 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0] :: \text{real}^3$
and $v5 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } 2 \ 0, \text{rat5 } 1 \ (-1), \text{rat5 } 0 \ 0] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4, v5\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v5\} \vee e = \text{convex hull } \{v2, v4\} \vee e = \text{convex hull } \{v3, v4\} \vee e = \text{convex hull } \{v3, v5\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0 rat5-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } (-1) \ (-1)$
unfolding assms
by $(\text{auto simp add: vector3-dot rat5-mul rat5-add})$
have $\text{finite}: \text{finite } \{v1, v2, v3, v4, v5\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$
qed

lemma *dodecahedron-facet3-edges*:

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) (1/2)] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0] :: \text{real}^3$
and $v4 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1/2) (1/2), \text{rat5 } (1/2) (1/2)] :: \text{real}^3$

and $v5 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (1/2) \ (1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } (-3) \ 1, \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4, v5\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v5\} \vee e = \text{convex hull } \{v3, v4\} \vee e = \text{convex hull } \{v4, v5\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0 rat5-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } 2 \ 0$
unfolding assms
by $(\text{auto simp add: vector3-dot rat5-mul rat5-add})$
have $\text{finite}: \text{finite } \{v1, v2, v3, v4, v5\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$
qed

lemma *dodecahedron-facet4-edges*:

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2)] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0] :: \text{real}^3$
and $v4 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2)] :: \text{real}^3$
and $v5 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } 3 \ (-1), \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4, v5\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v5\} \vee e = \text{convex hull } \{v3, v4\} \vee e = \text{convex hull } \{v4, v5\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0 rat5-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } (-2) \ 0$
unfolding assms
by $(\text{auto simp add: vector3-dot rat5-mul rat5-add})$
have $\text{finite}: \text{finite } \{v1, v2, v3, v4, v5\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$
qed

lemma *dodecahedron-facet5-edges*:

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0] :: \text{real}^3$
and $v4 \equiv \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0] :: \text{real}^3$
and $v5 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ (-1), \text{rat5 } (-3) \ 1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4, v5\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v4\} \vee e = \text{convex hull } \{v2, v3\} \vee e = \text{convex hull } \{v3, v5\} \vee e = \text{convex hull } \{v4, v5\})$

proof –

have $n\text{-nz}: n \neq 0$
 unfolding $n\text{-def}$ by (simp add: vector3-eq-0 rat5-eq-0)
 have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } 2 \ 0$
 unfolding assms
 by (auto simp add: vector3-dot rat5-mul rat5-add)
 have $\text{finite}: \text{finite } \{v1, v2, v3, v4, v5\}$
 using assms by blast
 show ?thesis
 by (facet-edges plane: plane $n\text{-nz}$: $n\text{-nz}$ defs: assms)

qed

lemma *dodecahedron-facet6-edges*:

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
 and $v2 \equiv \text{vector}[\text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
 and $v3 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0] :: \text{real}^3$
 and $v4 \equiv \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0] :: \text{real}^3$
 and $v5 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (1/2) \ (1/2)] :: \text{real}^3$
 and $n \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 1, \text{rat5 } (-3) \ 1] :: \text{real}^3$
 shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4, v5\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v4\} \vee e = \text{convex hull } \{v2, v3\} \vee e = \text{convex hull } \{v3, v5\} \vee e = \text{convex hull } \{v4, v5\})$

proof –

have $n\text{-nz}: n \neq 0$
 unfolding $n\text{-def}$ by (simp add: vector3-eq-0 rat5-eq-0)
 have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } (-2) \ 0$
 unfolding assms
 by (auto simp add: vector3-dot rat5-mul rat5-add)
 have $\text{finite}: \text{finite } \{v1, v2, v3, v4, v5\}$
 using assms by blast
 show ?thesis
 by (facet-edges plane: plane $n\text{-nz}$: $n\text{-nz}$ defs: assms)

qed

lemma *dodecahedron-facet7-edges*:

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
 and $v2 \equiv \text{vector}[\text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2)] :: \text{real}^3$
 and $v3 \equiv \text{vector}[\text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2)] :: \text{real}^3$
 and $v4 \equiv \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0] :: \text{real}^3$
 and $v5 \equiv \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0] :: \text{real}^3$
 and $n \equiv \text{vector}[\text{rat5 } (-2) \ 0, \text{rat5 } (-1) \ 1, \text{rat5 } 0 \ 0] :: \text{real}^3$
 shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4, v5\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v4\} \vee e = \text{convex hull } \{v1, v5\} \vee e = \text{convex hull } \{v2, v3\} \vee e = \text{convex hull } \{v2, v5\} \vee e = \text{convex hull } \{v3, v4\})$

proof –

have $n\text{-nz}: n \neq 0$
 unfolding $n\text{-def}$ by (simp add: vector3-eq-0 rat5-eq-0)
 have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } (-1) \ (-1)$
 unfolding assms

by (auto simp add: vector3-dot rat5-mul rat5-add)
 have finite: finite {v1, v2, v3, v4, v5}
 using assms by blast
 show ?thesis
 by (facet-edges plane: plane n-nz: n-nz defs: assms)
 qed

lemma dodecahedron-facet8-edges:

defines v1 $\equiv$ vector[rat5 (-1/2) (1/2), rat5 (1/2) (1/2), rat5 0 0] :: real³
 and v2 $\equiv$ vector[rat5 (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0] :: real³
 and v3 $\equiv$ vector[rat5 (-1) 0, rat5 1 0, rat5 (-1) 0] :: real³
 and v4 $\equiv$ vector[rat5 1 0, rat5 1 0, rat5 (-1) 0] :: real³
 and v5 $\equiv$ vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2)] :: real³
 and n $\equiv$ vector[rat5 0 0, rat5 1 (-1), rat5 3 (-1)] :: real³
 shows (e face-of (convex hull {v1, v2, v3, v4, v5}) $\wedge$ aff-dim e = 1) =
 (e = convex hull {v1, v2} $\vee$ e = convex hull {v1, v4} $\vee$ e = convex hull
 {v2, v3} $\vee$ e = convex hull {v3, v5} $\vee$ e = convex hull {v4, v5})
 proof -
 have n-nz: n $\neq$ 0
 unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
 have plane: $\bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } (-2) \ 0$
 unfolding assms
 by (auto simp add: vector3-dot rat5-mul rat5-add)
 have finite: finite {v1, v2, v3, v4, v5}
 using assms by blast
 show ?thesis
 by (facet-edges plane: plane n-nz: n-nz defs: assms)
 qed

lemma dodecahedron-facet9-edges:

defines v1 $\equiv$ vector[rat5 (-1/2) (1/2), rat5 (1/2) (1/2), rat5 0 0] :: real³
 and v2 $\equiv$ vector[rat5 (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0] :: real³
 and v3 $\equiv$ vector[rat5 (-1) 0, rat5 1 0, rat5 1 0] :: real³
 and v4 $\equiv$ vector[rat5 1 0, rat5 1 0, rat5 1 0] :: real³
 and v5 $\equiv$ vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (1/2) (1/2)] :: real³
 and n $\equiv$ vector[rat5 0 0, rat5 (-1) 1, rat5 3 (-1)] :: real³
 shows (e face-of (convex hull {v1, v2, v3, v4, v5}) $\wedge$ aff-dim e = 1) =
 (e = convex hull {v1, v2} $\vee$ e = convex hull {v1, v4} $\vee$ e = convex hull
 {v2, v3} $\vee$ e = convex hull {v3, v5} $\vee$ e = convex hull {v4, v5})
 proof -
 have n-nz: n $\neq$ 0
 unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
 have plane: $\bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } 2 \ 0$
 unfolding assms
 by (auto simp add: vector3-dot rat5-mul rat5-add)
 have finite: finite {v1, v2, v3, v4, v5}
 using assms by blast
 show ?thesis
 by (facet-edges plane: plane n-nz: n-nz defs: assms)

qed

lemma *dodecahedron-facet10-edges:*

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) (1/2), \text{rat5 } (1/2) (1/2), \text{rat5 } 0 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2)] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2)] :: \text{real}^3$
and $v4 \equiv \text{vector}[\text{rat5 } 1 0, \text{rat5 } 1 0, \text{rat5 } (-1) 0] :: \text{real}^3$
and $v5 \equiv \text{vector}[\text{rat5 } 1 0, \text{rat5 } 1 0, \text{rat5 } 1 0] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } 2 0, \text{rat5 } (-1) 1, \text{rat5 } 0 0] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4, v5\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v4\} \vee e = \text{convex hull } \{v1, v5\} \vee e = \text{convex hull } \{v2, v3\} \vee e = \text{convex hull } \{v2, v5\} \vee e = \text{convex hull } \{v3, v4\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0 rat5-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } 1 1$
unfolding assms
by $(\text{auto simp add: vector3-dot rat5-mul rat5-add})$
have $\text{finite}: \text{finite } \{v1, v2, v3, v4, v5\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$

qed

lemma *dodecahedron-facet11-edges:*

defines $v1 \equiv \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2)] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } 1 0, \text{rat5 } (-1) 0, \text{rat5 } 1 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } 1 0, \text{rat5 } 1 0, \text{rat5 } 1 0] :: \text{real}^3$
and $v4 \equiv \text{vector}[\text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2), \text{rat5 } (1/2) (1/2)] :: \text{real}^3$
and $v5 \equiv \text{vector}[\text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2), \text{rat5 } (1/2) (1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } (-3) 1, \text{rat5 } 0 0, \text{rat5 } 1 (-1)] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4, v5\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v5\} \vee e = \text{convex hull } \{v3, v4\} \vee e = \text{convex hull } \{v4, v5\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0 rat5-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } (-2) 0$
unfolding assms
by $(\text{auto simp add: vector3-dot rat5-mul rat5-add})$
have $\text{finite}: \text{finite } \{v1, v2, v3, v4, v5\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$

qed

lemma *dodecahedron-facet12-edges:*

defines $v1 \equiv \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2)] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } 1 0, \text{rat5 } (-1) 0, \text{rat5 } (-1) 0] :: \text{real}^3$

and $v3 \equiv \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0] :: \text{real}^3$
and $v4 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2)] :: \text{real}^3$
and $v5 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } 3 \ (-1), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ (-1)] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3, v4, v5\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v5\} \vee e = \text{convex hull } \{v3, v4\} \vee e = \text{convex hull } \{v4, v5\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0 rat5-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3, v4, v5\} \implies n \cdot x = \text{rat5 } 2 \ 0$
unfolding assms
by $(\text{auto simp add: vector3-dot rat5-mul rat5-add})$
have $\text{finite}: \text{finite } \{v1, v2, v3, v4, v5\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$
qed

lemmas $\text{dodecahedron-facet-edges} =$
 $\text{dodecahedron-facet1-edges dodecahedron-facet2-edges dodecahedron-facet3-edges}$
 $\text{dodecahedron-facet4-edges dodecahedron-facet5-edges dodecahedron-facet6-edges}$
 $\text{dodecahedron-facet7-edges dodecahedron-facet8-edges dodecahedron-facet9-edges}$
 $\text{dodecahedron-facet10-edges dodecahedron-facet11-edges dodecahedron-facet12-edges}$

lemma $\text{dodecahedron-edges-per-face:}$
assumes $f \text{ face-of std-dodecahedron}$
and $\text{aff-dim } f = 2$
shows $\text{card } \{e. e \text{ face-of std-dodecahedron} \wedge \text{aff-dim } e = 1 \wedge e \subseteq f\} = 5$
using $\text{conjI}[OF \text{ assms}]$ **unfolding** $\text{dodecahedron-facets face-edge-set-eq}[OF \text{ assms}(1)]$
apply $(\text{elim disjE forw-subst})$
unfolding $\text{dodecahedron-facet-edges set-cases card-5-iff}$
 $\text{segment-convex-hull[symmetric] closed-segment-eq}$
by $(\text{simp-all add: vector3-eq doubleton-eq-iff rat5-eq})$

lemma $\text{dodecahedron-edges:}$
 $e \text{ face-of std-dodecahedron} \wedge \text{aff-dim } e = 1 \iff$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2)],$
 $\text{vector}[\text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2)],$
 $\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2)],$
 $\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2)],$
 $\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2)],$
 $\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0],$
 $\text{vector}[\text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0]\} \vee$

```

e = convex hull {vector[rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 1 0, rat5 (-1) 0, rat5 (-1) 0]} ∨
e = convex hull {vector[rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 1 0, rat5 (-1) 0, rat5 1 0]} ∨
e = convex hull {vector[rat5 (-1/2) (1/2), rat5 (1/2) (1/2), rat5 0 0], vec-
tor[rat5 (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0]} ∨
e = convex hull {vector[rat5 (-1/2) (1/2), rat5 (1/2) (1/2), rat5 0 0], vec-
tor[rat5 1 0, rat5 1 0, rat5 (-1) 0]} ∨
e = convex hull {vector[rat5 (-1/2) (1/2), rat5 (1/2) (1/2), rat5 0 0], vec-
tor[rat5 1 0, rat5 1 0, rat5 1 0]} ∨
e = convex hull {vector[rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 (-1) 0, rat5 (-1) 0, rat5 (-1) 0]} ∨
e = convex hull {vector[rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0]} ∨
e = convex hull {vector[rat5 (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0], vec-
tor[rat5 (-1) 0, rat5 1 0, rat5 (-1) 0]} ∨
e = convex hull {vector[rat5 (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0], vec-
tor[rat5 (-1) 0, rat5 1 0, rat5 1 0]} ∨
e = convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1/2) (1/2)], vec-
tor[rat5 (1/2) (1/2), rat5 0 0, rat5 (1/2) (-1/2)]} ∨
e = convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1/2) (1/2)], vec-
tor[rat5 1 0, rat5 (-1) 0, rat5 1 0]} ∨
e = convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1/2) (1/2)], vec-
tor[rat5 1 0, rat5 1 0, rat5 1 0]} ∨
e = convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (1/2) (-1/2)], vec-
tor[rat5 1 0, rat5 (-1) 0, rat5 (-1) 0]} ∨
e = convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (1/2) (-1/2)], vec-
tor[rat5 1 0, rat5 1 0, rat5 (-1) 0]} ∨
e = convex hull {vector[rat5 (-1) 0, rat5 (-1) 0, rat5 (-1) 0], vector[rat5 0
0, rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2)]} ∨
e = convex hull {vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0], vector[rat5 0 0,
rat5 (1/2) (-1/2), rat5 (1/2) (1/2)]} ∨
e = convex hull {vector[rat5 (-1) 0, rat5 1 0, rat5 (-1) 0], vector[rat5 0 0,
rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2)]} ∨
e = convex hull {vector[rat5 (-1) 0, rat5 1 0, rat5 1 0], vector[rat5 0 0, rat5
(-1/2) (1/2), rat5 (1/2) (1/2)]} ∨
e = convex hull {vector[rat5 1 0, rat5 (-1) 0, rat5 (-1) 0], vector[rat5 0 0,
rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2)]} ∨
e = convex hull {vector[rat5 1 0, rat5 (-1) 0, rat5 1 0], vector[rat5 0 0, rat5
(1/2) (-1/2), rat5 (1/2) (1/2)]} ∨
e = convex hull {vector[rat5 1 0, rat5 1 0, rat5 (-1) 0], vector[rat5 0 0, rat5
(-1/2) (1/2), rat5 (-1/2) (-1/2)]} ∨
e = convex hull {vector[rat5 1 0, rat5 1 0, rat5 1 0], vector[rat5 0 0, rat5 (-1/2)
(1/2), rat5 (1/2) (1/2)]} ∨
e = convex hull {vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2)],
vector[rat5 0 0, rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2)]} ∨
e = convex hull {vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (1/2) (1/2)], vec-
tor[rat5 0 0, rat5 (1/2) (-1/2), rat5 (1/2) (1/2)]}
(is ?lhs  $\longleftrightarrow$  ?rhs)

```

```

proof
  assume ?lhs
  then obtain  $f$  where  $f$ -facts:  $f$  face-of std-dodecahedron  $\wedge$  aff-dim  $f = 2$   $e$  face-of
 $f \wedge$  aff-dim  $e = 1$ 
    using dodecahedron-polyhedron dodecahedron-fulldim edge-belongs-to-face
    by blast
  show ?rhs
    using  $f$ -facts unfolding dodecahedron-facets
    apply (elim disjE) apply (all ‹hyps subst-thin›)
    unfolding dodecahedron-facet-edges
    by linarith+
next
  assume ?rhs
  then show ?lhs
    using dodecahedron-facets dodecahedron-facet-edges
    by (smt (verit, ccfv-threshold) face-of-trans insertI1 insert-commute mem-Collect-eq)
qed

```

lemma dodecahedron-vertices:

```

shows ( $v$  face-of std-dodecahedron  $\wedge$  aff-dim  $v = 0$ ) =
  ( $v = \{ \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } 1 \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1) \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } 1 \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1) \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2), \text{rat5 } (1/2) \ (1/2)] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2)] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (1/2) \ (1/2)] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2)] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (-1/2) \ (1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (1/2) \ (-1/2), \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (1/2) \ (-1/2), \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2)] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1/2) \ (1/2)] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2)] \} \vee$ 
 $v = \{ \text{vector}[\text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } (1/2) \ (-1/2)] \}$ )
(is ?lhs = ?rhs)

```

```

proof
  assume *: ?lhs
  then obtain  $e$  where  $e$  face-of std-dodecahedron aff-dim  $e = 1$   $v$  face-of  $e$ 
    using dodecahedron-polyhedron dodecahedron-fulldim vertex-belongs-to-edge
    by blast
  show ?rhs
    using conjI[OF ‹ $e$  face-of std-dodecahedron› ‹aff-dim  $e = 1$ ›]
    unfolding dodecahedron-edges

```

```

    using vertices-of-edge[ $OF - \langle v \text{ face-of } e \rangle$ ] *
    by metis+
next
  assume ?rhs
  then show ?lhs
    apply (elim disjE forw-subst)
    by (smt (z3) aff-dim-sing extreme-point-of-segment face-of-singleton face-of-trans
        segment-convex-hull dodecahedron-edges)+
qed

```

```

lemma dodecahedron-edges-per-vertex:
  assumes  $v \text{ face-of std-dodecahedron}$ 
    and  $\text{aff-dim } v = 0$ 
  shows  $\text{card } \{e. e \text{ face-of std-dodecahedron} \wedge \text{aff-dim } e = 1 \wedge v \subseteq e\} = 3$ 
  unfolding vertex-edge-set-eq[ $OF \text{ assms}(1)$ ]
  using conjI[ $OF \text{ assms}$ ]
  unfolding dodecahedron-vertices conj-assoc[symmetric] dodecahedron-edges
  by (elim disjE forw-subst) edges-per-vertex+

```

8.3 Regularity

```

lemma dodecahedron-congruent-edges:
  assumes  $e1 \text{ face-of std-dodecahedron} \wedge \text{aff-dim } e1 = 1$ 
    and  $e2 \text{ face-of std-dodecahedron} \wedge \text{aff-dim } e2 = 1$ 
  shows  $e1 \text{ congruent } e2$ 
proof -
  let ?edges =
    {convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)]::real^3},
    vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)]},
    convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)],
    vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0]},
    convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)],
    vector[rat5 (-1) 0, rat5 1 0, rat5 1 0]},
    convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)],
    vector[rat5 (-1) 0, rat5 (-1) 0, rat5 (-1) 0]},
    convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)],
    vector[rat5 (-1) 0, rat5 1 0, rat5 (-1) 0]},
    convex hull {vector[rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2), rat5 0 0],
    vector[rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2), rat5 0 0]},
    convex hull {vector[rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2), rat5 0 0],
    vector[rat5 1 0, rat5 (-1) 0, rat5 (-1) 0]},
    convex hull {vector[rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2), rat5 0 0],
    vector[rat5 1 0, rat5 (-1) 0, rat5 1 0]},
    convex hull {vector[rat5 (-1/2) (1/2), rat5 (1/2) (1/2), rat5 0 0], vector[rat5
    (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0]},
    convex hull {vector[rat5 (-1/2) (1/2), rat5 (1/2) (1/2), rat5 0 0], vector[rat5
    1 0, rat5 1 0, rat5 (-1) 0]},
    convex hull {vector[rat5 (-1/2) (1/2), rat5 (1/2) (1/2), rat5 0 0], vector[rat5
    1 0, rat5 1 0, rat5 1 0]},

```

```

      convex hull {vector[rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 (-1) 0, rat5 (-1) 0, rat5 (-1) 0]},
      convex hull {vector[rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0]},
      convex hull {vector[rat5 (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0], vector[rat5
(-1) 0, rat5 1 0, rat5 (-1) 0]},
      convex hull {vector[rat5 (1/2) (-1/2), rat5 (1/2) (1/2), rat5 0 0], vector[rat5
(-1) 0, rat5 1 0, rat5 1 0]},
      convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5
(1/2) (1/2), rat5 0 0, rat5 (1/2) (-1/2)]},
      convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5
1 0, rat5 (-1) 0, rat5 1 0]},
      convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5
1 0, rat5 1 0, rat5 1 0]},
      convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (1/2) (-1/2)], vector[rat5
1 0, rat5 (-1) 0, rat5 (-1) 0]},
      convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (1/2) (-1/2)], vector[rat5
1 0, rat5 1 0, rat5 (-1) 0]},
      convex hull {vector[rat5 (-1) 0, rat5 (-1) 0, rat5 (-1) 0], vector[rat5 0 0,
rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2)]},
      convex hull {vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0], vector[rat5 0 0, rat5
(1/2) (-1/2), rat5 (1/2) (1/2)]},
      convex hull {vector[rat5 (-1) 0, rat5 1 0, rat5 (-1) 0], vector[rat5 0 0, rat5
(-1/2) (1/2), rat5 (-1/2) (-1/2)]},
      convex hull {vector[rat5 (-1) 0, rat5 1 0, rat5 1 0], vector[rat5 0 0, rat5
(-1/2) (1/2), rat5 (1/2) (1/2)]},
      convex hull {vector[rat5 1 0, rat5 (-1) 0, rat5 (-1) 0], vector[rat5 0 0, rat5
(1/2) (-1/2), rat5 (-1/2) (-1/2)]},
      convex hull {vector[rat5 1 0, rat5 (-1) 0, rat5 1 0], vector[rat5 0 0, rat5
(1/2) (-1/2), rat5 (1/2) (1/2)]},
      convex hull {vector[rat5 1 0, rat5 1 0, rat5 (-1) 0], vector[rat5 0 0, rat5
(-1/2) (1/2), rat5 (-1/2) (-1/2)]},
      convex hull {vector[rat5 1 0, rat5 1 0, rat5 1 0], vector[rat5 0 0, rat5 (-1/2)
(1/2), rat5 (1/2) (1/2)]},
      convex hull {vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2)],
vector[rat5 0 0, rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2)]},
      convex hull {vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (1/2) (1/2)], vector[rat5
0 0, rat5 (1/2) (-1/2), rat5 (1/2) (1/2)]}
    have e1 ∈ ?edges e2 ∈ ?edges
    using assms unfolding dodecahedron-edges
    by auto
  moreover have ∀ e ∈ ?edges.
    convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)],
vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)]} congruent e
    apply (simp only: ball-simps(7) Set.ball-empty simp-thms(21))
    unfolding conj-assoc[symmetric] segment-convex-hull[symmetric]
    apply (rule conjI)+
    apply (all ⟨rule congruent-segments⟩)
    unfolding dist-norm vector3-sub

```

apply (*simp-all add: norm-eq rat5-sub del: cancel-comm-monoid-add-class.diff-cancel diff-self*)
by (*simp-all add: vector3-dot rat5-mul rat5-add rat5-eq*)
ultimately show *?thesis*
using *congruent-set* **by** *meson*
qed

lemma *dodecahedron-congruent-faces:*

assumes *f1 face-of std-dodecahedron* $\wedge$ *aff-dim f1 = 2*

and *f2 face-of std-dodecahedron* $\wedge$ *aff-dim f2 = 2*

shows *f1 congruent f2*

proof –

let *?faces =*

$\{ \text{convex hull } \{ \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2)]::\text{real}^3, \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2)], \text{vector}[\text{rat5 } (1/2) (-1/2), \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } (-1) 0, \text{rat5 } (-1) 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } (-1) 0, \text{rat5 } 1 0] \},$
 $\text{convex hull } \{ \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2)], \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2)], \text{vector}[\text{rat5 } (1/2) (-1/2), \text{rat5 } (1/2) (1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } 1 0, \text{rat5 } (-1) 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } 1 0, \text{rat5 } 1 0] \},$
 $\text{convex hull } \{ \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2)], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } (-1) 0, \text{rat5 } 1 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } 1 0, \text{rat5 } 1 0], \text{vector}[\text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2), \text{rat5 } (1/2) (1/2)], \text{vector}[\text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2), \text{rat5 } (1/2) (1/2)] \},$
 $\text{convex hull } \{ \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2)], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } (-1) 0, \text{rat5 } (-1) 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } 1 0, \text{rat5 } (-1) 0], \text{vector}[\text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2), \text{rat5 } (-1/2) (-1/2)], \text{vector}[\text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2), \text{rat5 } (-1/2) (-1/2)] \},$
 $\text{convex hull } \{ \text{vector}[\text{rat5 } (-1/2) (1/2), \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (1/2) (-1/2), \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } (-1) 0, \text{rat5 } (-1) 0], \text{vector}[\text{rat5 } 1 0, \text{rat5 } (-1) 0, \text{rat5 } (-1) 0], \text{vector}[\text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2), \text{rat5 } (-1/2) (-1/2)] \},$
 $\text{convex hull } \{ \text{vector}[\text{rat5 } (-1/2) (1/2), \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (1/2) (-1/2), \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } (-1) 0, \text{rat5 } 1 0], \text{vector}[\text{rat5 } 1 0, \text{rat5 } (-1) 0, \text{rat5 } 1 0], \text{vector}[\text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2), \text{rat5 } (1/2) (1/2)] \},$
 $\text{convex hull } \{ \text{vector}[\text{rat5 } (-1/2) (1/2), \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2)], \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 0, \text{rat5 } (1/2) (-1/2)], \text{vector}[\text{rat5 } 1 0, \text{rat5 } (-1) 0, \text{rat5 } (-1) 0], \text{vector}[\text{rat5 } 1 0, \text{rat5 } (-1) 0, \text{rat5 } 1 0] \},$
 $\text{convex hull } \{ \text{vector}[\text{rat5 } (-1/2) (1/2), \text{rat5 } (1/2) (1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (1/2) (-1/2), \text{rat5 } (1/2) (1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } 1 0, \text{rat5 } (-1) 0], \text{vector}[\text{rat5 } 1 0, \text{rat5 } 1 0, \text{rat5 } (-1) 0], \text{vector}[\text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2), \text{rat5 } (-1/2) (-1/2)] \},$
 $\text{convex hull } \{ \text{vector}[\text{rat5 } (-1/2) (1/2), \text{rat5 } (1/2) (1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (1/2) (-1/2), \text{rat5 } (1/2) (1/2), \text{rat5 } 0 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } 1 0, \text{rat5 } 1 0], \text{vector}[\text{rat5 } 1 0, \text{rat5 } 1 0, \text{rat5 } 1 0], \text{vector}[\text{rat5 } 0 0, \text{rat5 } (-1/2) (1/2), \text{rat5 } (1/2) (1/2)] \},$

```

    convex hull {vector[rat5 (-1/2) (1/2), rat5 (1/2) (1/2), rat5 0 0], vector[rat5
(1/2) (1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5 (1/2) (1/2), rat5 0 0, rat5
(1/2) (-1/2)], vector[rat5 1 0, rat5 1 0, rat5 (-1) 0], vector[rat5 1 0, rat5 1 0,
rat5 1 0]},
    convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1/2) (1/2)], vector[rat5
1 0, rat5 (-1) 0, rat5 1 0], vector[rat5 1 0, rat5 1 0, rat5 1 0], vector[rat5 0 0,
rat5 (-1/2) (1/2), rat5 (1/2) (1/2)], vector[rat5 0 0, rat5 (1/2) (-1/2), rat5
(1/2) (1/2)]},
    convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (1/2) (-1/2)], vec-
tor[rat5 1 0, rat5 (-1) 0, rat5 (-1) 0], vector[rat5 1 0, rat5 1 0, rat5 (-1) 0],
vector[rat5 0 0, rat5 (-1/2) (1/2), rat5 (-1/2) (-1/2)], vector[rat5 0 0, rat5
(1/2) (-1/2), rat5 (-1/2) (-1/2)]}
    have f1 ∈ ?faces f2 ∈ ?faces
    using assms dodecahedron-facets
    by auto
    moreover have ∀ f ∈ ?faces.
    convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)], vec-
tor[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)], vector[rat5 (1/2) (-1/2),
rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 (-1) 0, rat5 (-1) 0, rat5 (-1) 0],
vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0]} congruent f
    apply (simp only: ball-simps(7) Set.ball-empty simp-thms(21))
    unfolding conj-assoc[symmetric]
    by (rule conjI)+ show-congruence+
    ultimately show ?thesis
    using congruent-set by meson
qed

```

```

lemma dodecahedron-facet1-equiangular:
  defines v1 ≡ vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1/2) (1/2)] :: real^3
  defines v2 ≡ vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2) (-1/2)] :: real^3
  defines v3 ≡ vector[rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2), rat5 0 0] :: real^3
  defines v4 ≡ vector[rat5 (-1) 0, rat5 (-1) 0, rat5 (-1) 0] :: real^3
  defines v5 ≡ vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0] :: real^3
  shows equiangular (convex hull {v1, v2, v3, v4, v5}) (3 * pi / 5)
  unfolding equiangular-def
proof (clarify)
  fix e1 e2 v
  assume *: e1 face-of convex hull {v1, v2, v3, v4, v5} aff-dim e1 = 1
    e2 face-of convex hull {v1, v2, v3, v4, v5} aff-dim e2 = 1
    e1 ≠ e2 v extreme-point-of e1 v extreme-point-of e2
  have arccos (1/4 - 1/4 * sqrt 5) = 3 * pi / 5
  proof (rule arccos-unique)
    show 0 ≤ 3 * pi / 5 by simp
    show 3 * pi / 5 ≤ pi by simp

    define t where t = 3 * pi / 5
    define c where c = cos t
    define s where s = sin t

```

have $\cos (3 * t) = \cos (2 * t + t)$ **by** *simp*
 also have $\dots = (2 * c^2 - 1) * c - 2 * s^2 * c$
 unfolding *cos-add c-def s-def*
 by (*simp add: cos-double-cos power2-eq-square sin-double*)
 finally have *cos-3t*: $\cos (3 * t) = 4 * c^3 - 3 * c$
 unfolding *sin-squared-eq c-def s-def* **by** *algebra*

have $\sin (3 * t) = \sin (2 * t + t)$ **by** *simp*
 also have $\dots = 2 * s * c^2 + (1 - 2 * s^2) * s$
 unfolding *sin-add c-def s-def*
 by (*simp add: cos-double-sin power2-eq-square sin-double*)
 finally have *sin-3t*: $\sin (3 * t) = 3 * s - 4 * s^3$
 unfolding *cos-squared-eq c-def s-def* **by** *algebra*

have $\cos (5 * t) = \cos (3 * t + 2 * t)$ **by** *simp*
 also have $\dots = \cos (3 * t) * \cos (2 * t) - \sin (3 * t) * \sin (2 * t)$
 using *cos-add* **by** *blast*
 also have $\dots = (4 * c^3 - 3 * c) * (2 * c^2 - 1) - (3 * s - 4 * s^3) * 2$
 $* s * c$
 by (*simp add: cos-3t sin-3t c-def s-def cos-double-cos sin-double*)
 also have $\dots = (4 * c^3 - 3 * c) * (2 * c^2 - 1) - 6 * s^2 * c + 8 * (s^2)^2 * c$
 by *algebra*
 also have $\dots = (4 * c^3 - 3 * c) * (2 * c^2 - 1) - 6 * (1 - c^2) * c + 8 * (1 - c^2)^2 * c$
 unfolding *sin-squared-eq c-def s-def* **by** *blast*
 finally have $\cos (5 * t) = 16 * c^5 - 20 * c^3 + 5 * c$
 by *algebra*
 moreover have $\cos (5 * t) = -1$
 unfolding *t-def* **by** *simp*
 ultimately have $16 * c^5 - 20 * c^3 + 5 * c + 1 = 0$
 by *simp*
 then have $(c + 1) * (4 * c^2 - 2 * c - 1)^2 = 0$
 by *algebra*
 then have $c + 1 = 0 \vee 4 * c^2 - 2 * c - 1 = 0$
 by *simp*

have $\exists !x. 0 \leq x \wedge x \leq \pi \wedge \cos x = -1$
 using *cos-total[of -1]* **by** *linarith*
 moreover have $0 \leq \pi \wedge \pi \leq \pi \wedge \cos \pi = -1$
 using *cos-pi* **by** *simp*
 moreover have $3 * \pi / 5 \neq \pi \wedge 0 \leq 3 * \pi / 5 \wedge 3 * \pi / 5 \leq \pi$
 by *auto*
 ultimately have $c \neq -1$
 unfolding *c-def t-def* **by** *blast*
 then have $c + 1 \neq 0$
 by *linarith*

then have $4 * c^2 + (-2) * c + (-1) = 0$

```

    using ⟨ $c + 1 = 0 \vee 4 * c^2 - 2 * c - 1 = 0$ ⟩
    by simp
  moreover have  $\text{discrim } 4 \ (-2) \ (-1) = 20$ 
    unfolding  $\text{discrim-def}$  by simp
  moreover have  $\text{sqrt } 20 = 2 * \text{sqrt } 5$ 
  by (metis numeral-times-numeral real-sqrt-four real-sqrt-mult semiring-norm(12,14))
  ultimately have  $(c = (2 + 2 * \text{sqrt } 5) / (2 * 4) \vee$ 
     $c = (2 - 2 * \text{sqrt } 5) / (2 * 4))$ 
    using  $\text{discriminant-iff}[of\ 4\ c\ -2\ -1]$ 
    by auto
  then have  $c = 1/4 + 1/4 * \text{sqrt } 5 \vee c = 1/4 - 1/4 * \text{sqrt } 5$ 
    by fastforce
  moreover have  $c < 0$ 
    unfolding  $c\text{-def}\ t\text{-def}$  using  $\text{cos-lt-zero-pi}[of\ 3 * \text{pi} / 5]$ 
    by simp
  ultimately show  $\text{cos } (3 * \text{pi} / 5) = 1 / 4 - 1 / 4 * \text{sqrt } 5$ 
    unfolding  $c\text{-def}\ t\text{-def}$ 
  by (metis add-nonneg-nonneg less-le not-le real-rat5-le-sqrt-cases zero-less-divide-1-iff
    zero-less-numeral)
qed
then have  $\text{arccos-3pi5: arccos } (\text{rat5 } (1/4) \ (- (1/4))) = 3 * \text{pi} / 5$  unfolding
rat5-def
  by (simp add: of-rat-divide of-rat-minus)
have  $\text{rat5 } 6 \ (-2) \geq 0$ 
  unfolding  $\text{rat5-0-1}(1)[\text{symmetric}]$   $\text{rat5-le}$ 
  by simp
have  $(v = v1 \wedge ((e1 = \text{convex hull } \{v1, v2\} \wedge e2 = \text{convex hull } \{v1, v5\}) \vee$ 
   $(e1 = \text{convex hull } \{v1, v5\} \wedge e2 = \text{convex hull } \{v1, v2\}))) \vee$ 
   $(v = v2 \wedge ((e1 = \text{convex hull } \{v2, v1\} \wedge e2 = \text{convex hull } \{v2, v4\}) \vee$ 
   $(e1 = \text{convex hull } \{v2, v4\} \wedge e2 = \text{convex hull } \{v2, v1\}))) \vee$ 
   $(v = v3 \wedge ((e1 = \text{convex hull } \{v3, v4\} \wedge e2 = \text{convex hull } \{v3, v5\}) \vee$ 
   $(e1 = \text{convex hull } \{v3, v5\} \wedge e2 = \text{convex hull } \{v3, v4\}))) \vee$ 
   $(v = v4 \wedge ((e1 = \text{convex hull } \{v4, v2\} \wedge e2 = \text{convex hull } \{v4, v3\}) \vee$ 
   $(e1 = \text{convex hull } \{v4, v3\} \wedge e2 = \text{convex hull } \{v4, v2\}))) \vee$ 
   $(v = v5 \wedge ((e1 = \text{convex hull } \{v5, v1\} \wedge e2 = \text{convex hull } \{v5, v3\}) \vee$ 
   $(e1 = \text{convex hull } \{v5, v3\} \wedge e2 = \text{convex hull } \{v5, v1\})))$ 
  using  $\text{conjI}[OF\ *(1,2)]$   $\text{conjI}[OF\ *(3,4)]$   $*(5-7)$  unfolding  $\text{assms dodecahedron-facet1-edges}$ 
  by (smt (verit) divide-cancel-right divide-eq-0-iff extreme-point-of-convex-hull-2
    insert-commute
     $\text{rat5-eq solids-vectors-to-rat5}(10)\ \text{vector3-neq}$ )
  moreover have  $\text{vangle } (v2 - v1) \ (v5 - v1) = 3 * \text{pi} / 5$ 
    vangle  $(v1 - v2) \ (v4 - v2) = 3 * \text{pi} / 5$ 
    vangle  $(v4 - v3) \ (v5 - v3) = 3 * \text{pi} / 5$ 
    vangle  $(v2 - v4) \ (v3 - v4) = 3 * \text{pi} / 5$ 
    vangle  $(v1 - v5) \ (v3 - v5) = 3 * \text{pi} / 5$ 
  unfolding  $\text{assms vector3-sub norm-eq-sqrt-inner vector3-dot vangle-def rat5-sub}$ 
  rat5-mul
  by (simp-all add: ⟨ $\text{rat5 } 6 \ (-2) \geq 0$ ⟩  $\text{arccos-3pi5 vector3-eq-0 rat5-sub rat5-mul}$ )

```

```

      rat5-add rat5-eq-0 rat5-div)
ultimately show ( $\exists a\ b. e1 = \text{convex hull } \{v, a\} \wedge e2 = \text{convex hull } \{v, b\} \wedge$ 
      vangle (a - v) (b - v) = 3 * pi / 5)
using vangle-commute
by metis
qed

lemma dodecahedron-equiangular:
assumes f face-of std-dodecahedron  $\wedge$  aff-dim f = 2
shows equiangular f (3 * pi / 5)
proof -
  define f1 where f1 = convex hull {(vector[rat5 (-1/2) (-1/2), rat5 0 0,
  rat5 (-1/2) (1/2)]::real^3), vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (1/2)
  (-1/2)], vector[rat5 (1/2) (-1/2), rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5
  (-1) 0, rat5 (-1) 0, rat5 (-1) 0], vector[rat5 (-1) 0, rat5 (-1) 0, rat5 1 0]}
  have f1 face-of std-dodecahedron  $\wedge$  aff-dim f1 = 2
  using f1-def dodecahedron-facets by blast
  then have f1 congruent f
  using assms dodecahedron-congruent-faces by blast
  then show ?thesis
  using equiangular-congruent dodecahedron-facet1-equiangular f1-def
  by blast
qed

end

```

9 Icosahedron

```

theory Icosahedron
imports
  Computation
begin

```

9.1 Definition

```

definition std-icosahedron :: (real, 3) vec set
where std-icosahedron  $\equiv$  convex hull
  {vector [0, 1,  $\varphi$ ], vector [0, 1,  $-\varphi$ ],
  vector [0, -1,  $\varphi$ ], vector [0, -1,  $-\varphi$ ],
  vector [1,  $\varphi$ , 0], vector [1,  $-\varphi$ , 0],
  vector [-1,  $\varphi$ , 0], vector [-1,  $-\varphi$ , 0],
  vector [ $\varphi$ , 0, 1], vector [ $-\varphi$ , 0, 1],
  vector [ $\varphi$ , 0, -1], vector [ $-\varphi$ , 0, -1]}

```

```

lemma std-icosahedron-eq:
shows std-icosahedron = convex hull
  { vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)],
  vector[rat5 0 0, rat5 1 0, rat5 (-1/2) (-1/2)],
  vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2) (1/2)],

```

```

vector[rat5 0 0, rat5 (-1) 0, rat5 (-1/2) (-1/2)],
vector[rat5 1 0, rat5 (1/2) (1/2), rat5 0 0],
vector[rat5 1 0, rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 (-1) 0, rat5 (1/2) (1/2), rat5 0 0],
vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0],
vector[rat5 (1/2) (1/2), rat5 0 0, rat5 1 0],
vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0],
vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1) 0],
vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0]]}
unfolding std-icosahedron-def rat5-def
by (simp add: add-divide-distrib of-rat-divide of-rat-minus)

```

```

lemma icosahedron-fulldim:
shows aff-dim std-icosahedron = 3
unfolding std-icosahedron-def
apply (rule polytope-3D)
by (simp add: vector3-sub vector3-cross vector3-dot vector3-eq-0,
      smt (verit, ccfv-threshold) real-sqrt-ge-0-iff)+

```

```

lemma icosahedron-polyhedron:
shows polyhedron std-icosahedron
by (simp add: std-icosahedron-def polytope-convex-hull polytope-imp-polyhedron)

```

9.2 Facets, edges, and vertices

```

lemma icosahedron-facets:
fixes f :: (real^3) set
shows f face-of std-icosahedron  $\wedge$  aff-dim f = 2  $\longleftrightarrow$ 
  f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0],
vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0], vector[rat5 (-1) 0, rat5 (-1/2)
(-1/2), rat5 0 0]}  $\vee$ 
  f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0],
vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0], vector[rat5 (-1) 0, rat5 (1/2)
(1/2), rat5 0 0]}  $\vee$ 
  f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0],
vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 0 0, rat5 (-1) 0,
rat5 (-1/2) (-1/2)]}  $\vee$ 
  f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0],
vector[rat5 (-1) 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 0 0, rat5 1 0, rat5
(-1/2) (-1/2)]}  $\vee$ 
  f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0],
vector[rat5 0 0, rat5 (-1) 0, rat5 (-1/2) (-1/2)], vector[rat5 0 0, rat5 1 0, rat5
(-1/2) (-1/2)]}  $\vee$ 
  f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0], vector[rat5
(-1) 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2)
(1/2)]}  $\vee$ 
  f = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0], vector[rat5
(-1) 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)]}
 $\vee$ 

```

$f = \text{convex hull} \{ \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2)], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2)] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2)] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2)] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2)], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2)] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2)] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2)] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2)] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2)] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2)] \}$
 $\vee$
 $f = \text{convex hull} \{ \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2)] \}$
unfolding std-icosahedron-eq
apply (simp only: compute-faces-2 finite-insert finite.emptyI)
apply (simp only: compute-faces-2-step-1 bex-empty simp-thms(31))
apply (simp only: compute-faces-2-step-2 bex-empty simp-thms(31))
apply (simp only: empty-iff ball-insert bex-empty bex-simps(5) ball-simps(5) simp-thms(21,31))
apply (simp only: vector3-sub vector3-cross vector3-eq-0 rat5-add rat5-sub rat5-mul rat5-eq)
apply (simp add: vector3-eq-0 vector3-dot rat5-add rat5-sub rat5-mul rat5-le rat5-eq rat5-eq-0 Let-def)
— very slow, ca 60s
apply (simp only: insert-commute)
by linarith

lemma *icosahedron-facet1-edges*:

```

defines v1  $\equiv$  vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0] :: real^3
and v2  $\equiv$  vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0] :: real^3
and v3  $\equiv$  vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0] :: real^3
and n  $\equiv$  vector[rat5 1 1, rat5 (-1) 1, rat5 0 0] :: real^3
shows (e face-of (convex hull {v1, v2, v3})  $\wedge$  aff-dim e = 1) =
      (e = convex hull {v1, v2}  $\vee$  e = convex hull {v1, v3}  $\vee$  e = convex hull
{v2, v3})
proof -
  have n-nz: n  $\neq$  0
  unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
  have plane:  $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } (-3) (-1)$ 
  unfolding v1-def v2-def v3-def n-def
  by (auto simp add: vector3-dot rat5-mul rat5-add)
  have finite: finite {v1, v2, v3}
  using assms by blast
  show ?thesis
  by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

```

lemma *icosahedron-facet2-edges*:

```

defines v1  $\equiv$  vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0] :: real^3
and v2  $\equiv$  vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0] :: real^3
and v3  $\equiv$  vector[rat5 (-1) 0, rat5 (1/2) (1/2), rat5 0 0] :: real^3
and n  $\equiv$  vector[rat5 (-1) (-1), rat5 (-1) 1, rat5 0 0] :: real^3
shows (e face-of (convex hull {v1, v2, v3})  $\wedge$  aff-dim e = 1) =
      (e = convex hull {v1, v2}  $\vee$  e = convex hull {v1, v3}  $\vee$  e = convex hull
{v2, v3})
proof -
  have n-nz: n  $\neq$  0
  unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
  have plane:  $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } 3 1$ 
  unfolding v1-def v2-def v3-def n-def
  by (auto simp add: vector3-dot rat5-mul rat5-add)
  have finite: finite {v1, v2, v3}
  using assms by blast
  show ?thesis
  by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

```

lemma *icosahedron-facet3-edges*:

```

defines v1  $\equiv$  vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0] :: real^3
and v2  $\equiv$  vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0] :: real^3
and v3  $\equiv$  vector[rat5 0 0, rat5 (-1) 0, rat5 (-1/2) (-1/2)] :: real^3
and n  $\equiv$  vector[rat5 2 0, rat5 2 0, rat5 2 0] :: real^3
shows (e face-of (convex hull {v1, v2, v3})  $\wedge$  aff-dim e = 1) =
      (e = convex hull {v1, v2}  $\vee$  e = convex hull {v1, v3}  $\vee$  e = convex hull
{v2, v3})

```

```

proof –
  have  $n\text{-nz}: n \neq 0$ 
    unfolding  $n\text{-def}$  by (simp add: vector3-eq-0 rat5-eq-0)
  have  $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } (-3) (-1)$ 
    unfolding  $v1\text{-def } v2\text{-def } v3\text{-def } n\text{-def}$ 
    by (auto simp add: vector3-dot rat5-mul rat5-add)
  have  $\text{finite}: \text{finite } \{v1, v2, v3\}$ 
    using assms by blast
  show ?thesis
    by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemma icosahedron-facet4-edges:
  defines  $v1 \equiv \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0\ 0, \text{rat5 } (-1)\ 0] :: \text{real}^3$ 
    and  $v2 \equiv \text{vector}[\text{rat5 } (-1)\ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0\ 0] :: \text{real}^3$ 
    and  $v3 \equiv \text{vector}[\text{rat5 } 0\ 0, \text{rat5 } 1\ 0, \text{rat5 } (-1/2) (-1/2)] :: \text{real}^3$ 
    and  $n \equiv \text{vector}[\text{rat5 } (-2)\ 0, \text{rat5 } 2\ 0, \text{rat5 } (-2)\ 0] :: \text{real}^3$ 
  shows  $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$ 
     $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$ 
proof –
  have  $n\text{-nz}: n \neq 0$ 
    unfolding  $n\text{-def}$  by (simp add: vector3-eq-0 rat5-eq-0)
  have  $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } 3\ 1$ 
    unfolding  $v1\text{-def } v2\text{-def } v3\text{-def } n\text{-def}$ 
    by (auto simp add: vector3-dot rat5-mul rat5-add)
  have  $\text{finite}: \text{finite } \{v1, v2, v3\}$ 
    using assms by blast
  show ?thesis
    by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemma icosahedron-facet5-edges:
  defines  $v1 \equiv \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0\ 0, \text{rat5 } (-1)\ 0] :: \text{real}^3$ 
    and  $v2 \equiv \text{vector}[\text{rat5 } 0\ 0, \text{rat5 } (-1)\ 0, \text{rat5 } (-1/2) (-1/2)] :: \text{real}^3$ 
    and  $v3 \equiv \text{vector}[\text{rat5 } 0\ 0, \text{rat5 } 1\ 0, \text{rat5 } (-1/2) (-1/2)] :: \text{real}^3$ 
    and  $n \equiv \text{vector}[\text{rat5 } (-1)\ 1, \text{rat5 } 0\ 0, \text{rat5 } 1\ 1] :: \text{real}^3$ 
  shows  $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$ 
     $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$ 
proof –
  have  $n\text{-nz}: n \neq 0$ 
    unfolding  $n\text{-def}$  by (simp add: vector3-eq-0 rat5-eq-0)
  have  $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } (-3) (-1)$ 
    unfolding  $v1\text{-def } v2\text{-def } v3\text{-def } n\text{-def}$ 
    by (auto simp add: vector3-dot rat5-mul rat5-add)
  have  $\text{finite}: \text{finite } \{v1, v2, v3\}$ 
    using assms by blast
  show ?thesis

```

by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemma icosahedron-facet6-edges:

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } (-2) \ 0, \text{rat5 } (-2) \ 0, \text{rat5 } 2 \ 0] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** (simp add: vector3-eq-0 rat5-eq-0)
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } 3 \ 1$
unfolding $v1\text{-def } v2\text{-def } v3\text{-def } n\text{-def}$
by (auto simp add: vector3-dot rat5-mul rat5-add)
have $\text{finite}: \text{finite } \{v1, v2, v3\}$
using assms **by** blast
show ?thesis
by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemma icosahedron-facet7-edges:

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } 2 \ 0, \text{rat5 } (-2) \ 0, \text{rat5 } (-2) \ 0] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** (simp add: vector3-eq-0 rat5-eq-0)
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } (-3) (-1)$
unfolding $v1\text{-def } v2\text{-def } v3\text{-def } n\text{-def}$
by (auto simp add: vector3-dot rat5-mul rat5-add)
have $\text{finite}: \text{finite } \{v1, v2, v3\}$
using assms **by** blast
show ?thesis
by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

lemma icosahedron-facet8-edges:

defines $v1 \equiv \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2)] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } 1 (-1), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$

```

      (e = convex hull {v1, v2} ∨ e = convex hull {v1, v3} ∨ e = convex hull
{v2, v3})
proof -
  have n-nz: n ≠ 0
    unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
  have plane:  $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } 3 \ 1$ 
    unfolding v1-def v2-def v3-def n-def
    by (auto simp add: vector3-dot rat5-mul rat5-add)
  have finite: finite {v1, v2, v3}
    using assms by blast
  show ?thesis
    by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

```

lemma icosahedron-facet9-edges:

```

defines v1 ≡ vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1) 0] :: real^3
  and v2 ≡ vector[rat5 (1/2) (1/2), rat5 0 0, rat5 1 0] :: real^3
  and v3 ≡ vector[rat5 1 0, rat5 (-1/2) (-1/2), rat5 0 0] :: real^3
  and n ≡ vector[rat5 1 1, rat5 1 (-1), rat5 0 0] :: real^3
shows (e face-of (convex hull {v1, v2, v3}) ∧ aff-dim e = 1) =
      (e = convex hull {v1, v2} ∨ e = convex hull {v1, v3} ∨ e = convex hull
{v2, v3})
proof -
  have n-nz: n ≠ 0
    unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
  have plane:  $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } 3 \ 1$ 
    unfolding v1-def v2-def v3-def n-def
    by (auto simp add: vector3-dot rat5-mul rat5-add)
  have finite: finite {v1, v2, v3}
    using assms by blast
  show ?thesis
    by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

```

lemma icosahedron-facet10-edges:

```

defines v1 ≡ vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1) 0] :: real^3
  and v2 ≡ vector[rat5 (1/2) (1/2), rat5 0 0, rat5 1 0] :: real^3
  and v3 ≡ vector[rat5 1 0, rat5 (1/2) (1/2), rat5 0 0] :: real^3
  and n ≡ vector[rat5 (-1) (-1), rat5 1 (-1), rat5 0 0] :: real^3
shows (e face-of (convex hull {v1, v2, v3}) ∧ aff-dim e = 1) =
      (e = convex hull {v1, v2} ∨ e = convex hull {v1, v3} ∨ e = convex hull
{v2, v3})
proof -
  have n-nz: n ≠ 0
    unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
  have plane:  $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } (-3) \ (-1)$ 
    unfolding v1-def v2-def v3-def n-def
    by (auto simp add: vector3-dot rat5-mul rat5-add)
  have finite: finite {v1, v2, v3}

```

using *assms* by *blast*
 show ?thesis
 by (facet-edges plane: plane n-nz: n-nz defs: *assms*)
 qed

lemma *icosahedron-facet11-edges*:

defines $v1 \equiv \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0\ 0, \text{rat5 } (-1) 0] :: \text{real}^3$
 and $v2 \equiv \text{vector}[\text{rat5 } 1\ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0\ 0] :: \text{real}^3$
 and $v3 \equiv \text{vector}[\text{rat5 } 0\ 0, \text{rat5 } (-1) 0, \text{rat5 } (-1/2) (-1/2)] :: \text{real}^3$
 and $n \equiv \text{vector}[\text{rat5 } 2\ 0, \text{rat5 } (-2) 0, \text{rat5 } (-2) 0] :: \text{real}^3$
 shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
 proof -
 have $n\text{-nz}: n \neq 0$
 unfolding *n-def* by (simp add: *vector3-eq-0 rat5-eq-0*)
 have *plane*: $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } 3\ 1$
 unfolding *v1-def v2-def v3-def n-def*
 by (auto simp add: *vector3-dot rat5-mul rat5-add*)
 have *finite*: *finite* $\{v1, v2, v3\}$
 using *assms* by *blast*
 show ?thesis
 by (facet-edges plane: plane n-nz: n-nz defs: *assms*)
 qed

lemma *icosahedron-facet12-edges*:

defines $v1 \equiv \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0\ 0, \text{rat5 } (-1) 0] :: \text{real}^3$
 and $v2 \equiv \text{vector}[\text{rat5 } 1\ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0\ 0] :: \text{real}^3$
 and $v3 \equiv \text{vector}[\text{rat5 } 0\ 0, \text{rat5 } 1\ 0, \text{rat5 } (-1/2) (-1/2)] :: \text{real}^3$
 and $n \equiv \text{vector}[\text{rat5 } (-2) 0, \text{rat5 } (-2) 0, \text{rat5 } 2\ 0] :: \text{real}^3$
 shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
 proof -
 have $n\text{-nz}: n \neq 0$
 unfolding *n-def* by (simp add: *vector3-eq-0 rat5-eq-0*)
 have *plane*: $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } (-3) (-1)$
 unfolding *v1-def v2-def v3-def n-def*
 by (auto simp add: *vector3-dot rat5-mul rat5-add*)
 have *finite*: *finite* $\{v1, v2, v3\}$
 using *assms* by *blast*
 show ?thesis
 by (facet-edges plane: plane n-nz: n-nz defs: *assms*)
 qed

lemma *icosahedron-facet13-edges*:

defines $v1 \equiv \text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0\ 0, \text{rat5 } (-1) 0] :: \text{real}^3$
 and $v2 \equiv \text{vector}[\text{rat5 } 0\ 0, \text{rat5 } (-1) 0, \text{rat5 } (-1/2) (-1/2)] :: \text{real}^3$
 and $v3 \equiv \text{vector}[\text{rat5 } 0\ 0, \text{rat5 } 1\ 0, \text{rat5 } (-1/2) (-1/2)] :: \text{real}^3$

```

    and n ≡ vector[rat5 (-1) 1, rat5 0 0, rat5 (-1) (-1)] :: real^3
  shows (e face-of (convex hull {v1, v2, v3}) ∧ aff-dim e = 1) =
    (e = convex hull {v1, v2} ∨ e = convex hull {v1, v3} ∨ e = convex hull
{v2, v3})
proof -
  have n-nz: n ≠ 0
  unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
  have plane: ∧x. x ∈ {v1, v2, v3} ⇒ n · x = rat5 3 1
  unfolding v1-def v2-def v3-def n-def
  by (auto simp add: vector3-dot rat5-mul rat5-add)
  have finite: finite {v1, v2, v3}
  using assms by blast
  show ?thesis
  by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

```

lemma *icosahedron-facet14-edges:*

```

  defines v1 ≡ vector[rat5 (1/2) (1/2), rat5 0 0, rat5 1 0] :: real^3
    and v2 ≡ vector[rat5 1 0, rat5 (-1/2) (-1/2), rat5 0 0] :: real^3
    and v3 ≡ vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2) (1/2)] :: real^3
    and n ≡ vector[rat5 (-2) 0, rat5 2 0, rat5 (-2) 0] :: real^3
  shows (e face-of (convex hull {v1, v2, v3}) ∧ aff-dim e = 1) =
    (e = convex hull {v1, v2} ∨ e = convex hull {v1, v3} ∨ e = convex hull
{v2, v3})
proof -
  have n-nz: n ≠ 0
  unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
  have plane: ∧x. x ∈ {v1, v2, v3} ⇒ n · x = rat5 (-3) (-1)
  unfolding v1-def v2-def v3-def n-def
  by (auto simp add: vector3-dot rat5-mul rat5-add)
  have finite: finite {v1, v2, v3}
  using assms by blast
  show ?thesis
  by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

```

lemma *icosahedron-facet15-edges:*

```

  defines v1 ≡ vector[rat5 (1/2) (1/2), rat5 0 0, rat5 1 0] :: real^3
    and v2 ≡ vector[rat5 1 0, rat5 (1/2) (1/2), rat5 0 0] :: real^3
    and v3 ≡ vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)] :: real^3
    and n ≡ vector[rat5 2 0, rat5 2 0, rat5 2 0] :: real^3
  shows (e face-of (convex hull {v1, v2, v3}) ∧ aff-dim e = 1) =
    (e = convex hull {v1, v2} ∨ e = convex hull {v1, v3} ∨ e = convex hull
{v2, v3})
proof -
  have n-nz: n ≠ 0
  unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
  have plane: ∧x. x ∈ {v1, v2, v3} ⇒ n · x = rat5 3 1
  unfolding v1-def v2-def v3-def n-def

```

by (auto simp add: vector3-dot rat5-mul rat5-add)
 have finite: finite {v1, v2, v3}
 using assms by blast
 show ?thesis
 by (facet-edges plane: plane n-nz: n-nz defs: assms)
 qed

lemma icosahedron-facet16-edges:

defines v1 $\equiv$ vector[rat5 (1/2) (1/2), rat5 0 0, rat5 1 0] :: real³
 and v2 $\equiv$ vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2) (1/2)] :: real³
 and v3 $\equiv$ vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)] :: real³
 and n $\equiv$ vector[rat5 1 (-1), rat5 0 0, rat5 (-1) (-1)] :: real³
 shows (e face-of (convex hull {v1, v2, v3}) $\wedge$ aff-dim e = 1) =
 (e = convex hull {v1, v2} $\vee$ e = convex hull {v1, v3} $\vee$ e = convex hull
 {v2, v3})
 proof -
 have n-nz: n $\neq$ 0
 unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
 have plane: $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } (-3) (-1)$
 unfolding v1-def v2-def v3-def n-def
 by (auto simp add: vector3-dot rat5-mul rat5-add)
 have finite: finite {v1, v2, v3}
 using assms by blast
 show ?thesis
 by (facet-edges plane: plane n-nz: n-nz defs: assms)
 qed

lemma icosahedron-facet17-edges:

defines v1 $\equiv$ vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0] :: real³
 and v2 $\equiv$ vector[rat5 1 0, rat5 (-1/2) (-1/2), rat5 0 0] :: real³
 and v3 $\equiv$ vector[rat5 0 0, rat5 (-1) 0, rat5 (-1/2) (-1/2)] :: real³
 and n $\equiv$ vector[rat5 0 0, rat5 1 1, rat5 (-1) 1] :: real³
 shows (e face-of (convex hull {v1, v2, v3}) $\wedge$ aff-dim e = 1) =
 (e = convex hull {v1, v2} $\vee$ e = convex hull {v1, v3} $\vee$ e = convex hull
 {v2, v3})
 proof -
 have n-nz: n $\neq$ 0
 unfolding n-def by (simp add: vector3-eq-0 rat5-eq-0)
 have plane: $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } (-3) (-1)$
 unfolding v1-def v2-def v3-def n-def
 by (auto simp add: vector3-dot rat5-mul rat5-add)
 have finite: finite {v1, v2, v3}
 using assms by blast
 show ?thesis
 by (facet-edges plane: plane n-nz: n-nz defs: assms)
 qed

lemma icosahedron-facet18-edges:

defines v1 $\equiv$ vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0] :: real³

and $v2 \equiv \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1/2) \ (-1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1/2) \ (1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ (-1), \text{rat5 } (-1) \ 1] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0 rat5-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } 3 \ 1$
unfolding $v1\text{-def } v2\text{-def } v3\text{-def } n\text{-def}$
by $(\text{auto simp add: vector3-dot rat5-mul rat5-add})$
have $\text{finite}: \text{finite } \{v1, v2, v3\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$
qed

lemma *icosahedron-facet19-edges:*

defines $v1 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1/2) \ (-1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 1, \text{rat5 } 1 \ (-1)] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0 rat5-eq-0})$
have $\text{plane}: \bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } 3 \ 1$
unfolding $v1\text{-def } v2\text{-def } v3\text{-def } n\text{-def}$
by $(\text{auto simp add: vector3-dot rat5-mul rat5-add})$
have $\text{finite}: \text{finite } \{v1, v2, v3\}$
using assms **by** blast
show $?thesis$
by $(\text{facet-edges plane: plane } n\text{-nz: } n\text{-nz defs: assms})$
qed

lemma *icosahedron-facet20-edges:*

defines $v1 \equiv \text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v2 \equiv \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) \ (1/2), \text{rat5 } 0 \ 0] :: \text{real}^3$
and $v3 \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (1/2) \ (1/2)] :: \text{real}^3$
and $n \equiv \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ (-1), \text{rat5 } 1 \ (-1)] :: \text{real}^3$
shows $(e \text{ face-of } (\text{convex hull } \{v1, v2, v3\}) \wedge \text{aff-dim } e = 1) =$
 $(e = \text{convex hull } \{v1, v2\} \vee e = \text{convex hull } \{v1, v3\} \vee e = \text{convex hull } \{v2, v3\})$
proof –
have $n\text{-nz}: n \neq 0$
unfolding $n\text{-def}$ **by** $(\text{simp add: vector3-eq-0 rat5-eq-0})$

```

have plane:  $\bigwedge x. x \in \{v1, v2, v3\} \implies n \cdot x = \text{rat5 } (-3) (-1)$ 
  unfolding v1-def v2-def v3-def n-def
  by (auto simp add: vector3-dot rat5-mul rat5-add)
have finite: finite {v1, v2, v3}
  using assms by blast
show ?thesis
  by (facet-edges plane: plane n-nz: n-nz defs: assms)
qed

```

```

lemmas icosahedron-facet-edges =
  icosahedron-facet1-edges icosahedron-facet2-edges icosahedron-facet3-edges
  icosahedron-facet4-edges icosahedron-facet5-edges icosahedron-facet6-edges
  icosahedron-facet7-edges icosahedron-facet8-edges icosahedron-facet9-edges
  icosahedron-facet10-edges icosahedron-facet11-edges icosahedron-facet12-edges
  icosahedron-facet13-edges icosahedron-facet14-edges icosahedron-facet15-edges
  icosahedron-facet16-edges icosahedron-facet17-edges icosahedron-facet18-edges
  icosahedron-facet19-edges icosahedron-facet20-edges

```

```

lemma icosahedron-edges-per-face:
  assumes f face-of std-icosahedron
    and aff-dim f = 2
  shows card {e. e face-of std-icosahedron  $\wedge$  aff-dim e = 1  $\wedge$  e  $\subseteq$  f} = 3
  using conjI[OF assms] unfolding icosahedron-facets face-edge-set-eq[OF assms(1)]
  apply (elim disjE forw-subst)
  unfolding icosahedron-facet-edges set-cases card-3-iff2
    segment-convex-hull[symmetric] closed-segment-eq
  by (simp-all add: vector3-eq doubleton-eq-iff rat5-eq)

```

```

lemma icosahedron-edges:
  e face-of std-icosahedron  $\wedge$  aff-dim e = 1  $\longleftrightarrow$ 
  e = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0], vector[rat5
(-1/2) (-1/2), rat5 0 0, rat5 1 0]}  $\vee$ 
  e = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0], vector[rat5
(-1) 0, rat5 (-1/2) (-1/2), rat5 0 0]}  $\vee$ 
  e = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0], vector[rat5
(-1) 0, rat5 (1/2) (1/2), rat5 0 0]}  $\vee$ 
  e = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0], vector[rat5
0 0, rat5 (-1) 0, rat5 (-1/2) (-1/2)]}  $\vee$ 
  e = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0], vector[rat5
0 0, rat5 1 0, rat5 (-1/2) (-1/2)]}  $\vee$ 
  e = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0], vector[rat5
(-1) 0, rat5 (-1/2) (-1/2), rat5 0 0]}  $\vee$ 
  e = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0], vector[rat5
(-1) 0, rat5 (1/2) (1/2), rat5 0 0]}  $\vee$ 
  e = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0], vector[rat5 0
0, rat5 (-1) 0, rat5 (1/2) (1/2)]}  $\vee$ 
  e = convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0], vector[rat5 0
0, rat5 1 0, rat5 (1/2) (1/2)]}  $\vee$ 
  e = convex hull {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1) 0], vector[rat5

```

$(1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2), \text{rat5 } 0 \ 0], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (-1/2) (-1/2)], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (-1/2) (-1/2)]\} \vee$
 $e = \text{convex hull } \{\text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } (-1) \ 0, \text{rat5 } (1/2) (1/2)], \text{vector}[\text{rat5 } 0 \ 0, \text{rat5 } 1 \ 0, \text{rat5 } (1/2) (1/2)]\}$
(is ?lhs $\longleftrightarrow$?rhs)

proof

assume ?lhs

then obtain f **where** f -facts: f face-of std-icosahedron $\wedge$ aff-dim $f = 2$ e face-of $f \wedge$ aff-dim $e = 1$

using icosahedron-polyhedron icosahedron-fulldim edge-belongs-to-face

by blast

show ?rhs

```

    using f-facts unfolding icosahedron-facets
    apply (elim disjE) apply (all ⟨hyps subst-thin⟩)
    unfolding icosahedron-facet-edges
    by linarith+
next
  assume ?rhs
  then show ?lhs
    using icosahedron-facets icosahedron-facet-edges

  by (smt (verit, ccfv-threshold) face-of-trans insertI1 insert-commute mem-Collect-eq)
qed

```

lemma *icosahedron-vertices*:

```

shows (v face-of std-icosahedron ∧ aff-dim v = 0) =
  (v = {vector[rat5 0 0, rat5 1 0, rat5 (1/2) (1/2)]} ∨
    v = {vector[rat5 0 0, rat5 1 0, rat5 (-1/2) (-1/2)]} ∨
    v = {vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2) (1/2)]} ∨
    v = {vector[rat5 0 0, rat5 (-1) 0, rat5 (-1/2) (-1/2)]} ∨
    v = {vector[rat5 1 0, rat5 (1/2) (1/2), rat5 0 0]} ∨
    v = {vector[rat5 1 0, rat5 (-1/2) (-1/2), rat5 0 0]} ∨
    v = {vector[rat5 (-1) 0, rat5 (1/2) (1/2), rat5 0 0]} ∨
    v = {vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0]} ∨
    v = {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 1 0]} ∨
    v = {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0]} ∨
    v = {vector[rat5 (1/2) (1/2), rat5 0 0, rat5 (-1) 0]} ∨
    v = {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0]})
(is ?lhs = ?rhs)

```

proof

```

  assume *: ?lhs
  then obtain e where e face-of std-icosahedron aff-dim e = 1 v face-of e
    using icosahedron-polyhedron icosahedron-fulldim vertex-belongs-to-edge
  by blast
  show ?rhs
    using conjI[OF ⟨e face-of std-icosahedron⟩ ⟨aff-dim e = 1⟩]
    unfolding icosahedron-edges
    using vertices-of-edge[OF - ⟨v face-of e⟩] *
    by metis+
next
  assume ?rhs
  then show ?lhs
    apply (elim disjE forw-subst)
  by (smt (z3) aff-dim-sing extreme-point-of-segment face-of-singleton face-of-trans
    segment-convex-hull icosahedron-edges)+
qed

```

lemma *icosahedron-edges-per-vertex*:

```

assumes v face-of std-icosahedron
  and aff-dim v = 0
shows card {e. e face-of std-icosahedron ∧ aff-dim e = 1 ∧ v ⊆ e} = 5

```

unfolding *vertex-edge-set-eq*[*OF assms*(1)]
using *conjI*[*OF assms*]
unfolding *icosahedron-vertices conj-assoc*[*symmetric*] *icosahedron-edges*
by (*elim disjE* *forw-subst*) *edges-per-vertex*+

9.3 Regularity

lemma *icosahedron-congruent-edges*:

assumes *e1 face-of std-icosahedron* $\wedge$ *aff-dim e1 = 1*

and *e2 face-of std-icosahedron* $\wedge$ *aff-dim e2 = 1*

shows *e1 congruent e2*

proof –

let *?edges* =
 {*convex hull* {(*vector*[*rat5* (–1/2) (–1/2), *rat5* 0 0, *rat5* (–1) 0]::*real*³),
vector[*rat5* (–1/2) (–1/2), *rat5* 0 0, *rat5* 1 0]},
convex hull {*vector*[*rat5* (–1/2) (–1/2), *rat5* 0 0, *rat5* (–1) 0], *vector*[*rat5*
 (–1) 0, *rat5* (–1/2) (–1/2), *rat5* 0 0]},
convex hull {*vector*[*rat5* (–1/2) (–1/2), *rat5* 0 0, *rat5* (–1) 0], *vector*[*rat5*
 (–1) 0, *rat5* (1/2) (1/2), *rat5* 0 0]},
convex hull {*vector*[*rat5* (–1/2) (–1/2), *rat5* 0 0, *rat5* (–1) 0], *vector*[*rat5*
 0 0, *rat5* (–1) 0, *rat5* (–1/2) (–1/2)]},
convex hull {*vector*[*rat5* (–1/2) (–1/2), *rat5* 0 0, *rat5* (–1) 0], *vector*[*rat5*
 0 0, *rat5* 1 0, *rat5* (–1/2) (–1/2)]},
convex hull {*vector*[*rat5* (–1/2) (–1/2), *rat5* 0 0, *rat5* 1 0], *vector*[*rat5* (–1)
 0, *rat5* (–1/2) (–1/2), *rat5* 0 0]},
convex hull {*vector*[*rat5* (–1/2) (–1/2), *rat5* 0 0, *rat5* 1 0], *vector*[*rat5* (–1)
 0, *rat5* (1/2) (1/2), *rat5* 0 0]},
convex hull {*vector*[*rat5* (–1/2) (–1/2), *rat5* 0 0, *rat5* 1 0], *vector*[*rat5* 0 0,
rat5 (–1) 0, *rat5* (1/2) (1/2)]},
convex hull {*vector*[*rat5* (–1/2) (–1/2), *rat5* 0 0, *rat5* 1 0], *vector*[*rat5* 0 0,
rat5 1 0, *rat5* (1/2) (1/2)]},
convex hull {*vector*[*rat5* (1/2) (1/2), *rat5* 0 0, *rat5* (–1) 0], *vector*[*rat5*
 (1/2) (1/2), *rat5* 0 0, *rat5* 1 0]},
convex hull {*vector*[*rat5* (1/2) (1/2), *rat5* 0 0, *rat5* (–1) 0], *vector*[*rat5* 1
 0, *rat5* (–1/2) (–1/2), *rat5* 0 0]},
convex hull {*vector*[*rat5* (1/2) (1/2), *rat5* 0 0, *rat5* (–1) 0], *vector*[*rat5* 1
 0, *rat5* (1/2) (1/2), *rat5* 0 0]},
convex hull {*vector*[*rat5* (1/2) (1/2), *rat5* 0 0, *rat5* (–1) 0], *vector*[*rat5* 0
 0, *rat5* (–1) 0, *rat5* (–1/2) (–1/2)]},
convex hull {*vector*[*rat5* (1/2) (1/2), *rat5* 0 0, *rat5* (–1) 0], *vector*[*rat5* 0
 0, *rat5* 1 0, *rat5* (–1/2) (–1/2)]},
convex hull {*vector*[*rat5* (1/2) (1/2), *rat5* 0 0, *rat5* 1 0], *vector*[*rat5* 1 0,
rat5 (–1/2) (–1/2), *rat5* 0 0]},
convex hull {*vector*[*rat5* (1/2) (1/2), *rat5* 0 0, *rat5* 1 0], *vector*[*rat5* 1 0,
rat5 (1/2) (1/2), *rat5* 0 0]},
convex hull {*vector*[*rat5* (1/2) (1/2), *rat5* 0 0, *rat5* 1 0], *vector*[*rat5* 0 0,
rat5 (–1) 0, *rat5* (1/2) (1/2)]},
convex hull {*vector*[*rat5* (1/2) (1/2), *rat5* 0 0, *rat5* 1 0], *vector*[*rat5* 0 0,
rat5 1 0, *rat5* (1/2) (1/2)]},

```

      convex hull {vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5
1 0, rat5 (-1/2) (-1/2), rat5 0 0]},
      convex hull {vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5
0 0, rat5 (-1) 0, rat5 (-1/2) (-1/2)]},
      convex hull {vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5
0 0, rat5 (-1) 0, rat5 (1/2) (1/2)]},
      convex hull {vector[rat5 (-1) 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 1
0, rat5 (1/2) (1/2), rat5 0 0]},
      convex hull {vector[rat5 (-1) 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 0
0, rat5 1 0, rat5 (-1/2) (-1/2)]},
      convex hull {vector[rat5 (-1) 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 0
0, rat5 1 0, rat5 (1/2) (1/2)]},
      convex hull {vector[rat5 1 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 0 0,
rat5 (-1) 0, rat5 (-1/2) (-1/2)]},
      convex hull {vector[rat5 1 0, rat5 (-1/2) (-1/2), rat5 0 0], vector[rat5 0 0,
rat5 (-1) 0, rat5 (1/2) (1/2)]},
      convex hull {vector[rat5 1 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 0 0,
rat5 1 0, rat5 (-1/2) (-1/2)]},
      convex hull {vector[rat5 1 0, rat5 (1/2) (1/2), rat5 0 0], vector[rat5 0 0,
rat5 1 0, rat5 (1/2) (1/2)]},
      convex hull {vector[rat5 0 0, rat5 (-1) 0, rat5 (-1/2) (-1/2)], vector[rat5
0 0, rat5 1 0, rat5 (-1/2) (-1/2)]},
      convex hull {vector[rat5 0 0, rat5 (-1) 0, rat5 (1/2) (1/2)], vector[rat5 0
0, rat5 1 0, rat5 (1/2) (1/2)]}
    have e1 ∈ ?edges e2 ∈ ?edges
    using assms unfolding icosahedron-edges
    by auto
  moreover have ∀ e ∈ ?edges.
    convex hull {vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0], vector[rat5
(-1/2) (-1/2), rat5 0 0, rat5 1 0]} congruent e
    apply (simp only: ball-simps(7) Set.ball-empty simp-thms(21))
    unfolding conj-assoc[symmetric] segment-convex-hull[symmetric]
    apply (rule conjI)+
    apply (all ‹rule congruent-segments›)
    unfolding dist-norm vector3-sub
    apply (simp-all add: norm-eq rat5-sub del: cancel-comm-monoid-add-class.diff-cancel
diff-self)
    by (simp-all add: vector3-dot rat5-mul rat5-add rat5-eq)
  ultimately show ?thesis
    using congruent-set by meson
qed

```

lemma icosahedron-congruent-faces:

```

  assumes f1 face-of std-icosahedron ∧ aff-dim f1 = 2
    and f2 face-of std-icosahedron ∧ aff-dim f2 = 2
  shows f1 congruent f2
proof -
  let ?faces =
    {convex hull {(vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0]::real^3),

```

[illegible]

```

have f1 ∈ ?faces f2 ∈ ?faces
  using assms icosahedron-facets
  by auto
moreover have ∀ f ∈ ?faces.
  convex hull {(vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0]::real^3),
vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0], vector[rat5 (-1) 0, rat5 (-1/2)
(-1/2), rat5 0 0]} congruent f
  apply (simp only: ball-simps(7) Set.ball-empty simp-thms(21))
  unfolding conj-assoc[symmetric]
  by (rule conjI)+ show-congruence+
ultimately show ?thesis
  using congruent-set by meson
qed

lemma icosahedron-facet1-equiangular:
  defines v1 ≡ vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 (-1) 0] :: real^3
  defines v2 ≡ vector[rat5 (-1/2) (-1/2), rat5 0 0, rat5 1 0] :: real^3
  defines v3 ≡ vector[rat5 (-1) 0, rat5 (-1/2) (-1/2), rat5 0 0] :: real^3
  shows equiangular (convex hull {v1, v2, v3}) (pi / 3)
  unfolding equiangular-def
proof (clarify)
  fix e1 e2 v
  assume *: e1 face-of convex hull {v1, v2, v3} aff-dim e1 = 1
           e2 face-of convex hull {v1, v2, v3} aff-dim e2 = 1
           e1 ≠ e2 v extreme-point-of e1 v extreme-point-of e2
  have rat5 4 0 ≥ 0 rat5 (1/2) 0 = 1/2
    unfolding rat5-def
    by (simp-all add: of-rat-divide)
  have (v = v1 ∧ ((e1 = convex hull {v1, v2} ∧ e2 = convex hull {v1, v3}) ∨
    (e1 = convex hull {v1, v3} ∧ e2 = convex hull {v1, v2}))) ∨
    (v = v2 ∧ ((e1 = convex hull {v2, v1} ∧ e2 = convex hull {v2, v3}) ∨
    (e1 = convex hull {v2, v3} ∧ e2 = convex hull {v2, v1}))) ∨
    (v = v3 ∧ ((e1 = convex hull {v3, v1} ∧ e2 = convex hull {v3, v2}) ∨
    (e1 = convex hull {v3, v2} ∧ e2 = convex hull {v3, v1})))
    using conjI[OF *(1,2)] conjI[OF *(3,4)] *(5-7) unfolding assms icosahedron-facet1-edges
    by (smt (verit, best) extreme-point-of-convex-hull-2 insert-commute vector3-eq)
  moreover have vangle (v2 - v1) (v3 - v1) = pi / 3
    vangle (v1 - v2) (v3 - v2) = pi / 3
    vangle (v1 - v3) (v2 - v3) = pi / 3
  unfolding assms vector3-sub norm-eq-sqrt-inner vector3-dot vangle-def rat5-sub
  rat5-mul
  by (simp-all add: vector3-eq-0 rat5-sub rat5-mul rat5-add rat5-eq-0 rat5-div
    ⟨rat5 4 0 ≥ 0⟩ ⟨rat5 (1/2) 0 = 1/2⟩)
  ultimately show (∃ a b. e1 = convex hull {v, a} ∧ e2 = convex hull {v, b} ∧
    vangle (a - v) (b - v) = pi / 3)
    using vangle-commute
    by metis
qed

```

```

lemma icosahedron-equiangular:
  assumes  $f$  face-of std-icosahedron  $\wedge$  aff-dim  $f = 2$ 
  shows equiangular  $f$   $(\pi / 3)$ 
proof –
  define  $f1$  where  $f1 = \text{convex hull } \{(\text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0\ 0, \text{rat5 } (-1) 0]::\text{real}^3), \text{vector}[\text{rat5 } (-1/2) (-1/2), \text{rat5 } 0\ 0, \text{rat5 } 1\ 0], \text{vector}[\text{rat5 } (-1) 0, \text{rat5 } (-1/2) (-1/2), \text{rat5 } 0\ 0]\}$ 
  have  $f1$  face-of std-icosahedron  $\wedge$  aff-dim  $f1 = 2$ 
  using f1-def icosahedron-facets by blast
  then have  $f1$  congruent  $f$ 
  using assms icosahedron-congruent-faces by blast
  then show ?thesis
  using equiangular-congruent icosahedron-facet1-equiangular f1-def
  by blast
qed

end

```

10 Main result

```

theory Platonic-Solids
imports
  Euler-Polyhedron-Formula.Euler-Formula
  Polytope-More
  Tetrahedron
  Cube
  Octahedron
  Dodecahedron
  Icosahedron
begin

```

10.1 Possible Schläfli symbols

```

lemma platonic-solid-limits:
  fixes  $p :: (\text{real}^3)$  set
  assumes is-polytope: polytope  $p$ 
  and aff-dim-p: aff-dim  $p = 3$ 
  and edges-per-face:  $\forall f. f$  face-of  $p \wedge$  aff-dim  $f = 2 \longrightarrow$ 
     $\text{card } \{e. e$  face-of  $p \wedge$  aff-dim  $e = 1 \wedge e \subseteq f\} = m$ 
  and edges-per-vert:  $\forall v. v$  face-of  $p \wedge$  aff-dim  $v = 0 \longrightarrow$ 
     $\text{card } \{e. e$  face-of  $p \wedge$  aff-dim  $e = 1 \wedge v \subseteq e\} = n$ 
  shows  $(m = 3 \wedge n = 3) \vee$ 
     $(m = 4 \wedge n = 3) \vee$ 
     $(m = 3 \wedge n = 4) \vee$ 
     $(m = 5 \wedge n = 3) \vee$ 
     $(m = 3 \wedge n = 5)$ 
proof –
  let  $?V = \text{card } \{v. v$  face-of  $p \wedge$  aff-dim  $v = 0\}$ 

```

```

let ?E = card {e. e face-of p ∧ aff-dim e = 1}
let ?F = card {f. f face-of p ∧ aff-dim f = 2}
have ?V + ?F - ?E = 2
  using Euler-relation[OF is-polytope aff-dim-p] by fastforce
then have ?V + ?F = ?E + 2
  by linarith
moreover have ?V ≥ 4
  using polytope-vertex-lower-bound[OF is-polytope] aff-dim-p
  by fastforce
moreover have ?F ≥ 4
proof -
  have f facet-of p ⟷ f face-of p ∧ aff-dim f = 2 for f
    unfolding facet-of-def aff-dim-p by auto
  then show ?thesis
    using polytope-facet-lower-bound[OF is-polytope] aff-dim-p by presburger
qed
ultimately have ?E ≥ 6
  by linarith
have ⋀ f. f face-of p ⟹ aff-dim f = 2 ⟹ card {e. e face-of p ∧ aff-dim e = 1
  ∧ e ⊆ f} ≥ 3
proof -
  fix f
  assume *: f face-of p aff-dim f = 2
  then have f-polytope: polytope f
    using face-of-polytope-polytope is-polytope
    by blast
  have e facet-of f ⟷ e face-of p ∧ aff-dim e = 1 ∧ e ⊆ f for e
    using face-of-face[OF *(1)] unfolding facet-of-def *(2)
    by auto
  then show card {e. e face-of p ∧ aff-dim e = 1 ∧ e ⊆ f} ≥ 3
    using polytope-facet-lower-bound[OF f-polytope] *(2)
    by presburger
qed
then have m ≥ 3
  using edges-per-face
  by (metis (mono-tags, lifting) Collect-empty-eq One-nat-def
    ‹?F ≥ 4› add-Suc-shift card.empty linorder-not-le
    numeral-Bit0 one-add-one plus-nat.add-0 zero-less-Suc)

have m * ?F = 2 * ?E
  using edges-per-face facet-ridge-relation[OF is-polytope] unfolding aff-dim-p
  by simp
moreover have n * ?V = 2 * ?E
  using edge-vertex-relation[OF is-polytope edges-per-vert] .
moreover have m * (n * ?V) + n * (m * ?F) = m * n * (?E + 2)
  using ‹?V + ?F = ?E + 2› by algebra
ultimately have E-m-n-relation: 2 * ?E * (m + n) = m * n * (?E + 2)
  by algebra

```

```

have  $n \neq 0$ 
  using  $\langle ?E \geq 6 \rangle$   $E$ - $m$ - $n$ -relation  $\langle 3 \leq m \rangle$   $rel$ - $simps$ (28)
  by fastforce
moreover have  $n \neq 1$ 
  using  $\langle ?F \geq 4 \rangle$   $\langle ?V + ?F = ?E + 2 \rangle$   $\langle n * ?V = 2 * ?E \rangle$  by force
moreover have  $n \neq 2$ 
  using  $E$ - $m$ - $n$ -relation  $\langle m * ?F = 2 * ?E \rangle$   $\langle ?E \geq 6 \rangle$   $\langle ?F \geq 4 \rangle$  by force
ultimately have  $n \geq 3$ 
  by auto

have  $?E * (m * n) < ?E * (2 * (m + n))$ 
  using  $\langle 3 \leq m \rangle$   $\langle n \neq 0 \rangle$   $E$ - $m$ - $n$ -relation
  by (simp add: algebra-simps)
then have  $m * n < 2 * (m + n)$ 
  by auto
then have  $int\ m * int\ n - 2 * int\ m - 2 * int\ n < 0$ 
  by (smt (verit, ccfv-SIG) diff-add-cancel diff-less-0-iff-less group-cancel.add2
mult-2
    of-nat-add of-nat-less-iff of-nat-mult)
then have  $int\ m * (int\ n - 2) - 2 * int\ n < 0$ 
  by (simp add: right-diff-distrib')
then have  $int\ m * (int\ n - 2) - 2 * (int\ n - 2) < 4$ 
  by simp
then have  $int$ - $m$ - $n$ -ineq:  $(int\ m - 2) * (int\ n - 2) < 4$ 
  using  $int$ - $distrib$ (3) by presburger
have  $(m - 2) * (int\ n - 2) < 4$ 
  using  $int$ - $m$ - $n$ -ineq  $\langle 3 \leq m \rangle$  by simp
then have  $int\ ((m - 2) * (n - 2)) < int\ 4$ 
  by (simp add: of-nat-diff-if)
then have  $m$ - $n$ -ineq:  $(m - 2) * (n - 2) < 4$ 
  by linarith

have  $m - 2 \geq 1$   $n - 2 \geq 1$ 
  using  $\langle 3 \leq m \rangle$   $\langle 3 \leq n \rangle$  by auto

have  $m \leq 5$ 
  proof (rule ccontr)
    assume  $\neg(m \leq 5)$ 
    then have *:  $m - 2 \geq 4$  by simp
    have  $(m - 2) * (n - 2) \geq 4$ 
      using  $mult$ - $le$ - $mono[OF * \langle n - 2 \geq 1 \rangle]$  by auto
    then show False
      using  $m$ - $n$ -ineq by linarith
  qed
have  $n \leq 5$ 
  proof (rule ccontr)
    assume  $\neg(n \leq 5)$ 
    then have *:  $n - 2 \geq 4$  by simp
    have  $(m - 2) * (n - 2) \geq 4$ 

```

using *mult-le-mono*[*OF* $\langle m - 2 \geq 1 \rangle *$] by *auto*
 then show *False*
 using *m-n-ineq* by *linarith*
 qed

have $m = 3 \vee m = 4 \vee m = 5$
 using $\langle 3 \leq m \rangle \langle m \leq 5 \rangle$ by *linarith*
 moreover have $n = 3 \vee n = 4 \vee n = 5$
 using $\langle 3 \leq n \rangle \langle n \leq 5 \rangle$ by *linarith*
 ultimately show *?thesis*
 using *m-n-ineq* by *presburger*
 qed

10.2 Exactly five Platonic solids

theorem *platonic-solids-full*:

shows $(\exists p::(\text{real}^3) \text{ set. polytope } p \wedge \text{aff-dim } p = 3 \wedge$
 $(\forall f. f \text{ face-of } p \wedge \text{aff-dim } f = 2 \longrightarrow$
 $\text{card } \{e. e \text{ face-of } p \wedge \text{aff-dim } e = 1 \wedge e \subseteq f\} = m) \wedge$
 $(\forall v. v \text{ face-of } p \wedge \text{aff-dim } v = 0 \longrightarrow$
 $\text{card } \{e. e \text{ face-of } p \wedge \text{aff-dim } e = 1 \wedge v \subseteq e\} = n) \wedge$
 $(\forall f1 f2. f1 \text{ face-of } p \wedge \text{aff-dim } f1 = 2 \wedge$
 $f2 \text{ face-of } p \wedge \text{aff-dim } f2 = 2 \longrightarrow f1 \text{ congruent } f2) \wedge$
 $(\forall e1 e2. e1 \text{ face-of } p \wedge \text{aff-dim } e1 = 1 \wedge$
 $e2 \text{ face-of } p \wedge \text{aff-dim } e2 = 1 \longrightarrow e1 \text{ congruent } e2) \wedge$
 $(\exists t. \forall f. f \text{ face-of } p \wedge \text{aff-dim } f = 2 \longrightarrow \text{equiangular } f t)) =$
 $((m = 3 \wedge n = 3) \vee$
 $(m = 4 \wedge n = 3) \vee$
 $(m = 3 \wedge n = 4) \vee$
 $(m = 5 \wedge n = 3) \vee$
 $(m = 3 \wedge n = 5))$
 (is *?lhs* = *?rhs*)

proof

assume *?lhs*
 then show *?rhs*
 using *platonic-solid-limits* by *auto*

next

assume *?rhs*
 then consider

$(m3-n3) \ m = 3 \wedge n = 3 \mid$
 $(m3-n4) \ m = 3 \wedge n = 4 \mid$
 $(m3-n5) \ m = 3 \wedge n = 5 \mid$
 $(m4-n3) \ m = 4 \wedge n = 3 \mid$
 $(m5-n3) \ m = 5 \wedge n = 3$

by *fast*

then show *?lhs*

proof (*cases*)

case *m3-n3*

moreover have *polytope std-tetrahedron*

```

    by (simp add: polytope-convex-hull std-tetrahedron-def)
  ultimately show ?thesis
  using tetrahedron-fulldim tetrahedron-edges-per-face tetrahedron-edges-per-vertex
    tetrahedron-congruent-edges tetrahedron-congruent-faces tetrahedron-equiangular
    by blast
next
case m3-n4
moreover have polytope std-octahedron
  by (simp add: polytope-convex-hull std-octahedron-def)
ultimately show ?thesis
using octahedron-fulldim octahedron-edges-per-face octahedron-edges-per-vertex
  octahedron-congruent-edges octahedron-congruent-faces octahedron-equiangular
  by blast
next
case m3-n5
moreover have polytope std-icosahedron
  by (simp add: polytope-convex-hull std-icosahedron-def)
ultimately show ?thesis
using icosahedron-fulldim icosahedron-edges-per-face icosahedron-edges-per-vertex
  icosahedron-congruent-edges icosahedron-congruent-faces icosahedron-equiangular
  by blast
next
case m4-n3
moreover have polytope std-cube
  by (simp add: polytope-convex-hull std-cube-def)
ultimately show ?thesis
using cube-fulldim cube-edges-per-face cube-edges-per-vertex
  cube-congruent-edges cube-congruent-faces cube-equiangular
  by blast
next
case m5-n3
moreover have polytope std-dodecahedron
  by (simp add: polytope-convex-hull std-dodecahedron-def)
ultimately show ?thesis
using dodecahedron-fulldim dodecahedron-edges-per-face dodecahedron-edges-per-vertex
  dodecahedron-congruent-edges dodecahedron-congruent-faces dodecahedron-equiangular
  by blast
qed
qed

```

This is the original top-level statement of the HOL Light formalization, kept here for parity. It omits the regularity condition on p .

theorem *platonic-solids*:

shows $(\exists p::(\text{real}^3) \text{ set. } \text{polytope } p \wedge \text{aff-dim } p = 3 \wedge$
 $(\forall f. f \text{ face-of } p \wedge \text{aff-dim } f = 2 \longrightarrow$
 $\text{card } \{e. e \text{ face-of } p \wedge \text{aff-dim } e = 1 \wedge e \subseteq f\} = m) \wedge$
 $(\forall v. v \text{ face-of } p \wedge \text{aff-dim } v = 0 \longrightarrow$
 $\text{card } \{e. e \text{ face-of } p \wedge \text{aff-dim } e = 1 \wedge v \subseteq e\} = n)) =$
 $((m = 3 \wedge n = 3) \vee$

```

      (m = 4 ∧ n = 3) ∨
      (m = 3 ∧ n = 4) ∨
      (m = 5 ∧ n = 3) ∨
      (m = 3 ∧ n = 5))
(is ?lhs = ?rhs)
proof
  assume ?lhs
  then show ?rhs
    using platonic-solid-limits by auto
next
  assume ?rhs
  then show ?lhs
    unfolding platonic-solids-full[symmetric]
    by meson
qed

end

```

References

- [1] J. Harrison. Proof of the Platonic solids theorem, accessed 2025. HOL Light repository.