

Notes on Gödel's and Scott's Variants of the Ontological Argument (Isabelle/HOL dataset)

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Abstract

Experimental studies with Isabelle/HOL on Kurt Gödel's modal ontological argument and Dana Scott's variant of it are presented. They implicitly answer some questions that the authors have received over the last decade(s). In addition, some new results are reported.

Our contribution is explained in full detail in [2]. This document presents the corresponding Isabelle/HOL dataset (which is only slightly modified to meet AFP requirements).

Contents

1	Introduction	2
2	Interactive and automated theorem proving	3
2.1	SurjectiveCantor.thy (Figure 2 of [2])	3
3	Mechanization of higher-order modal logic (HOML)	4
3.1	HOMLinHOL.thy (Figure 3 of [2])	4
3.2	TestsHOML.thy (Figure 4 of [2])	6
3.3	ModalFilter.thy (Figure 5 of [2])	7
4	Gödels ontological argument – 1970 manuscript	8
4.1	GoedelVariantHOML1.thy (Figure 6 of [2])	8
4.2	GoedelVariantHOML2.thy (Figure 7 of [2])	9
4.3	GoedelVariantHOML3.thy (Figure 8 of [2])	11
4.4	ThereIsNoEvil1.thy (Figure 10 of [2])	13
4.5	ThereIsNoEvil2.thy (Figure 11 of [2])	13
5	Scotts variant	14
5.1	ScottVariantHOML.thy (Figure 12 of [2])	14
5.2	HOMLinHOLonlyK.thy (slight variation of Figure 3 of [2]) . .	15
5.3	ScottVariantHOMLinK.thy (Figure 13 of [2])	17
5.4	ScottVariantHOMLAndersonQuant.thy (Figure 15 of [2]) . .	17

6 Appendix	19
6.1 GoedelVariantHOML1poss.thy (Figure 16 of [2])	19
6.2 GoedelVariantHOML2poss.thy (Figure 17 of [2])	20
6.3 GoedelVariantHOML3poss.thy (Figure 18 of [2])	22
6.4 ScottVariantHOMLposs.thy (Figure 19 of [2])	24
6.5 EvilDerivable.thy (Figure 20 of [2])	25
7 Further Appendices	27
7.1 GoedelVariantHOML1AndersonQuan.thy	27
7.2 GoedelVariantHOML2AndersonQuan.thy	28
7.3 GoedelVariantHOML3AndersonQuan.thy	30
7.4 HOMLinHOLonlyS4.thy (slight variation of Figure 3 of [2]) .	32
7.5 TestsHOMLinS4.thy	33
7.6 GoedelVariantHOML1inS4.thy	35
7.7 GoedelVariantHOML2inS4.thy	35
7.8 GoedelVariantHOML2possInS4.thy	36
7.9 GoedelVariantHOML3inS4.thy	37
7.10 GoedelVariantHOML3possInS4.thy	38
7.11 ScottVariantHOMLinS4.thy	39
7.12 ScottVariantHOMLpossInS4.thy	40

1 Introduction

The Isabelle/HOL dataset associated with [2] is presented. Compared to previous work on Gödel’s modal ontological argument as published in the Archive of Formal Proofs (AFP) [1, 4, 3], our dataset addresses several relevant and in some cases novel aspects, which are combined here in a single publication, including:

1. For the first time, Gödel’s original manuscript [5] has been formalized as closely as possible.
2. The inconsistency of the postulates in Gödel’s original manuscript is explained in detail.
3. Two different ways of eliminating this inconsistency are presented, one of which is novel.
4. Scott’s variant [6] of Gödel’s original manuscript is presented and compared with Gödel’s original variant.
5. In addition to logics S5 and K, the above variants are also studied for logic S4.

6. The above variants are tested for various combinations of of possibilist and actualist quantifiers for individuals (which has not been done systematically in previous publications).
7. Concepts of evil are examined and the derivability of evil is critically questioned.

The purpose of this AFP publication is essentially twofold. One motivation is to make the data sources associated with [2] available in a sustainable, well-maintained way. The other motivation is to support university education in higher-order modal logic by providing a small dataset for reuse that illustrates a systematically explored philosophical argument, emphasizing in particular different notions of quantification.

Compared to [2], the Isabelle sources presented here have been slightly modified to meet some AFP requirements. This concerns the commenting out of calls to sledghammer (to reduce computational resources) and some minor reformatting (e.g. insertion of new lines). The formalization code itself remains unchanged.

2 Interactive and automated theorem proving

2.1 SurjectiveCantor.thy (Figure 2 of [2])

The surjective Cantor theorem is used in [2] to illustrate some aspects of interactive and automated theorem proving in Isabelle/HOL as relevant for the paper. To keep the provided data material complete wrt. [2], we include these data sources also here.

```
theory SurjectiveCantor imports Main
begin

Surjective Cantor theorem: traditional interactive proof
theorem SurjectiveCantor:  $\neg(\exists G. \forall F::'a \Rightarrow \text{bool}. \exists X::'a. G X = F)$ 
proof
  assume 1:  $\exists G. \forall F::'a \Rightarrow \text{bool}. \exists X::'a. G X = F$ 
  obtain g::'a $\Rightarrow('a \Rightarrow \text{bool})$  where 2:  $\forall F. \exists X. g X = F$  using 1 by auto
  let ?F =  $\lambda X. \neg g X X$ 
  have 3:  $\exists Y. g Y = ?F$  using 2 by metis
  obtain a::'a where 4:  $g a = ?F$  using 3 by auto
  have 5:  $g a a = ?F a$  using 4 by metis
  have 6:  $g a a = (\neg g a a)$  using 5 by auto
  show False using 6 by auto
qed
```

Avoiding proof by contradiction (Fuenmayor & Benzmüller, 2021)

```
theorem SurjectiveCantor':  $\neg(\exists G. \forall F::'a \Rightarrow \text{bool}. \exists X::'a. G X = F)$ 
proof –
```

```

{fix g :: 'a⇒('a⇒bool)
  have 1: ∀ X.∃ Y.(¬g X Y) = (¬g Y Y) by auto
  have 2: ∀ X.∃ Y.(¬g X Y) = ((λZ.¬g Z Z) Y) using 1 by auto
  have 3: ∃ F.∀ X.∃ Y.(¬g X Y) = (F Y) using 2 by auto
  have 4: ∃ F.∀ X.¬(∀ Y.(g X Y) = (F Y)) using 3 by auto
  have ∃ F.∀ X.¬(g X = F) using 4 by metis
}
hence 5: ∀ G.∃ F::'a⇒bool.∀ X::'a.¬(G X = F) by auto
have 6: ¬(∃ G.∀ F::'a⇒bool.∃ X::'a. G X = F) using 5 by auto
thus ?thesis . — done, avoiding proof by contradiction
qed

```

Surjective Cantor theorem: automated proof by some internal/external theorem provers

```

theorem SurjectiveCantor'': ¬(∃ G.∀ F::'a⇒bool.∃ X::'a. G X = F)
nitpick[expect=none] — no counterexample found
— sledgehammer — most internal provers give up
— sledgehammer[remote_leo2 remote_leo3] — proof found: leo provers succeed
oops

```

Surjective Cantor theorem (wrong formalization attempt): the types are crucial

```

theorem SurjectiveCantor'''': ¬(∃ G.∀ F::'b.∃ X::'a. G X = F)
nitpick — counterexample found for card 'a = 1 and card 'b = 1: G=(λx. (a1
:= b1)
nitpick[satisfy, expect=genuine] — model found for card 'a = 1 and card 'b = 2
nitpick[card 'a=2, card 'b=3, expect=none] — no counterexample found
oops
end

```

3 Mechanization of higher-order modal logic (HOML)

3.1 HOMLinHOL.thy (Figure 3 of [2])

Shallow embedding of higher-order modal logic (HOML) in the classical higher-order logic (HOL) of Isabelle/HOL utilizing the logic-pluralistic LogiKEY methodology. Here logic S5 is introduced.

```

theory HOMLinHOL imports Main
begin

```

— Global parameters setting for the model finder nitpick and the parser; unimport for the reader

```

nitpick-params[user-axioms,expect=genuine,show-all,format=2,max-genuine=3]
declare[[syntax-ambiguity-warning=false]]

```

— Type i is associated with possible worlds and type e with entities
typedecl *i* — Possible worlds

typeddecl e — Individuals/entities
type-synonym $\sigma = i \Rightarrow \text{bool}$ — World-lifted propositions
type-synonym $\tau = e \Rightarrow \sigma$ — Modal properties

consts $R :: i \Rightarrow i \Rightarrow \text{bool}$ (-r-) — Accessibility relation between worlds

axiomatization where
 $Rrefl: \forall x. xrx \text{ and}$
 $Rsymm: \forall x y. xry \longrightarrow yrx \text{ and}$
 $Rtrans: \forall x y z. xry \wedge yrz \longrightarrow xrz$

— Logical connectives (operating on truth-sets)
abbreviation $Mbot :: \sigma (\perp)$ **where** $\perp \equiv \lambda w. \text{False}$
abbreviation $Mtop :: \sigma (\top)$ **where** $\top \equiv \lambda w. \text{True}$
abbreviation $Mneg :: \sigma \Rightarrow \sigma (\neg)$ [52]53 **where** $\neg\varphi \equiv \lambda w. \neg(\varphi w)$
abbreviation $Mand :: \sigma \Rightarrow \sigma \Rightarrow \sigma (\text{infixl} \wedge 50)$ **where** $\varphi \wedge \psi \equiv \lambda w. \varphi w \wedge \psi w$
abbreviation $Mor :: \sigma \Rightarrow \sigma \Rightarrow \sigma (\text{infixl} \vee 49)$ **where** $\varphi \vee \psi \equiv \lambda w. \varphi w \vee \psi w$
abbreviation $Mimp :: \sigma \Rightarrow \sigma \Rightarrow \sigma (\text{infixr} \supset 48)$ **where** $\varphi \supset \psi \equiv \lambda w. \varphi w \longrightarrow \psi w$
abbreviation $Mequiv :: \sigma \Rightarrow \sigma \Rightarrow \sigma (\text{infixl} \leftrightarrow 47)$ **where** $\varphi \leftrightarrow \psi \equiv \lambda w. \varphi w \longleftrightarrow \psi w$
abbreviation $Mbox :: \sigma \Rightarrow \sigma (\Box)$ [54]55 **where** $\Box \varphi \equiv \lambda w. \forall v. w \mathbf{r} v \longrightarrow \varphi v$
abbreviation $Mdia :: \sigma \Rightarrow \sigma (\Diamond)$ [54]55 **where** $\Diamond \varphi \equiv \lambda w. \exists v. w \mathbf{r} v \wedge \varphi v$
abbreviation $Mprimeq :: 'a \Rightarrow 'a \Rightarrow \sigma (=)$ **where** $x = y \equiv \lambda w. x = y$
abbreviation $Mprimneg :: 'a \Rightarrow 'a \Rightarrow \sigma (\neq)$ **where** $x \neq y \equiv \lambda w. x \neq y$
abbreviation $Mnegpred :: \tau \Rightarrow \tau (\sim)$ **where** $\sim \Phi \equiv \lambda x. \lambda w. \neg \Phi x w$
abbreviation $Mconpred :: \tau \Rightarrow \tau \Rightarrow \tau (\text{infixl} . 50)$ **where** $\Phi . \Psi \equiv \lambda x. \lambda w. \Phi x w \wedge \Psi x w$
abbreviation $Mexclor :: \sigma \Rightarrow \sigma \Rightarrow \sigma (\text{infixl} \vee^e 49)$ **where** $\varphi \vee^e \psi \equiv (\varphi \vee \psi) \wedge \neg(\varphi \wedge \psi)$

— Possibilist quantifiers (polymorphic)
abbreviation $Mallposs :: ('a \Rightarrow \sigma) \Rightarrow \sigma (\forall)$ **where** $\forall \Phi \equiv \lambda w. \forall x. \Phi x w$
abbreviation $Mallpossb$ (**binder** \forall) [8]9 **where** $\forall x. \varphi(x) \equiv \forall \varphi$
abbreviation $Mexiposs :: ('a \Rightarrow \sigma) \Rightarrow \sigma (\exists)$ **where** $\exists \Phi \equiv \lambda w. \exists x. \Phi x w$
abbreviation $Mexipossb$ (**binder** \exists) [8]9 **where** $\exists x. \varphi(x) \equiv \exists \varphi$

— Actualist quantifiers (for individuals/entities)
consts $existsAt :: e \Rightarrow \sigma (@)$
abbreviation $Mallact :: (e \Rightarrow \sigma) \Rightarrow \sigma (\forall^E)$ **where** $\forall^E \Phi \equiv \lambda w. \forall x. x @ w \longrightarrow \Phi x w$
abbreviation $Mallactb$ (**binder** \forall^E) [8]9 **where** $\forall^E x. \varphi(x) \equiv \forall^E \varphi$
abbreviation $Mexiact :: (e \Rightarrow \sigma) \Rightarrow \sigma (\exists^E)$ **where** $\exists^E \Phi \equiv \lambda w. \exists x. x @ w \wedge \Phi x w$
abbreviation $Mexiactb$ (**binder** \exists^E) [8]9 **where** $\exists^E x. \varphi(x) \equiv \exists^E \varphi$

— Leibniz equality (polymorphic)
abbreviation $Mleibeq :: 'a \Rightarrow 'a \Rightarrow \sigma (\equiv)$ **where** $x \equiv y \equiv \forall P. P x \supset P y$

— Meta-logical predicate for global validity
abbreviation $Mvalid :: \sigma \Rightarrow \text{bool} ([\cdot])$ **where** $[\psi] \equiv \forall w. \psi w$

end

3.2 TestsHOML.thy (Figure 4 of [2])

Tests and verifications of properties for the embedding of HOML (S5) in HOL.

```
theory TestsHOML imports HOMLinHOL
begin
```

— Test for S5 modal logic

```
lemma axM: [ $\Box\varphi \supset \varphi$ ] using Rrefl by blast
lemma axD: [ $\Box\varphi \supset \Diamond\varphi$ ] using Rrefl by blast
lemma axB: [ $\varphi \supset \Box\Diamond\varphi$ ] using Rsymm by blast
lemma ax4: [ $\Box\varphi \supset \Box\Box\varphi$ ] using Rtrans by blast
lemma ax5: [ $\Diamond\varphi \supset \Box\Diamond\varphi$ ] using Rsymm Rtrans by blast
```

— Test for Barcan and converse Barcan formula:

```
lemma BarcanAct: [ $(\forall^E x. \Box(\varphi x)) \supset \Box(\forall^E x. (\varphi x))$ ]
nitpick[expect=genuine] oops — Countermodel found
lemma ConvBarcanAct: [ $\Box(\forall^E x. (\varphi x)) \supset (\forall^E x. \Box(\varphi x))$ ]
nitpick[expect=genuine] oops — Countermodel found
lemma BarcanPoss: [ $(\forall x. \Box(\varphi x)) \supset \Box(\forall x. \varphi x)$ ] by blast
lemma ConvBarcanPoss: [ $\Box(\forall x. (\varphi x)) \supset (\forall x. \Box(\varphi x))$ ] by blast
```

— A simple Hilbert system for classical propositional logic is derived

```
lemma Hilbert-A1: [ $A \supset (B \supset A)$ ] by blast
lemma Hilbert-A2: [ $(A \supset (B \supset C)) \supset ((A \supset B) \supset (A \supset C))$ ] by blast
lemma Hilbert-MP: assumes [ $A$ ] and [ $A \supset B$ ] shows [ $B$ ] using assms by blast
```

— We have a polymorphic possibilist quantifier for which existential import holds

```
lemma Quant-1: assumes [ $A$ ] shows [ $\forall x::'a. A$ ] using assms by auto
```

— Existential import holds for possibilist quantifiers

```
lemma ExImPossibilist1: [ $\exists x::e. x = x$ ] by blast
lemma ExImPossibilist2: [ $\exists x::e. x \equiv x$ ] by blast
lemma ExImPossibilist3: [ $\exists x::e. x = t$ ] by blast
lemma ExImPossibilist4: [ $\exists x::'a. x \equiv t::'a$ ] by blast
lemma ExImPossibilist: [ $\exists x::'a. \top$ ] by blast
```

— We have an actualist quantifier for individuals for which existential import does not hold

```
lemma Quant-2: assumes [ $A$ ] shows [ $\forall^E x::e. A$ ] using assms by auto
```

— Existential import does not hold for our actualist quantifiers (for individuals)

```
lemma ExImActualist1: [ $\exists^E x::e. x = x$ ]
nitpick[card=1,expect=genuine] oops — Countermodel found
lemma ExImActualist2: [ $\exists^E x::e. x \equiv x$ ]
nitpick[card=1,expect=genuine] oops — Countermodel found
lemma ExImActualist3: [ $\exists^E x::e. x = t$ ]
nitpick[card=1,expect=genuine] oops — Countermodel found
lemma ExImActualist: [ $\exists^E x::e. \top$ ]
```

```

nitpick[card=1,expect=genuine] oops — Countermodel found

— Properties of the embedded primitive equality, which coincides with Leibniz
equality
lemma EqRefl:  $\lfloor x = x \rfloor$  by blast
lemma EqSym:  $\lfloor (x = y) \leftrightarrow (y = x) \rfloor$  by blast
lemma EqTrans:  $\lfloor ((x = y) \wedge (y = z)) \supset (x = z) \rfloor$  by blast
lemma EQCong:  $\lfloor (x = y) \supset ((\varphi x) = (\varphi y)) \rfloor$  by blast
lemma EQFuncExt:  $\lfloor (\varphi = \psi) \supset (\forall x. ((\varphi x) = (\psi x))) \rfloor$  by blast
lemma EQBoolExt1:  $\lfloor (\varphi = \psi) \supset (\varphi \leftrightarrow \psi) \rfloor$  by blast
lemma EQBoolExt2:  $\lfloor (\varphi \leftrightarrow \psi) \supset (\varphi = \psi) \rfloor$ 
nitpick[card=2] oops — Countermodel found
lemma EQBoolExt3:  $\lfloor (\varphi \leftrightarrow \psi) \rfloor \longrightarrow \lfloor (\varphi = \psi) \rfloor$  by blast
lemma EqPrimLeib:  $\lfloor (x = y) \leftrightarrow (x \equiv y) \rfloor$  by auto

— Comprehension is natively supported in HOL (due to lambda-abstraction)
lemma Comprehension1:  $\lfloor \exists \varphi. \forall x. (\varphi x) \leftrightarrow A \rfloor$  by force
lemma Comprehension2:  $\lfloor \exists \varphi. \forall x. (\varphi x) \leftrightarrow (A x) \rfloor$  by force
lemma Comprehension3:  $\lfloor \exists \varphi. \forall x y. (\varphi x y) \leftrightarrow (A x y) \rfloor$  by force

— Modal collapse does not hold
lemma ModalCollapse:  $\lfloor \forall \varphi. \varphi \supset \Box \varphi \rfloor$ 
nitpick[card=2,expect=genuine] oops — Countermodel found

— Empty property and self-difference
lemma TruePropertyAndSelfIdentity:  $\lfloor (\lambda x::e. \top) = (\lambda x. x = x) \rfloor$  by blast
lemma EmptyPropertyAndSelfDifference:  $\lfloor (\lambda x::e. \perp) = (\lambda x. x \neq x) \rfloor$  by blast
lemma EmptyProperty2:  $\lfloor \exists x. \varphi x \rfloor \longrightarrow \lfloor \varphi \neq (\lambda x::e. \perp) \rfloor$  by blast
lemma EmptyProperty3:  $\lfloor \exists^E x. \varphi x \rfloor \longrightarrow \lfloor \varphi \neq (\lambda x::e. \perp) \rfloor$  by blast
lemma EmptyProperty4:  $\lfloor \varphi \neq (\lambda x::e. \perp) \rfloor \longrightarrow \lfloor \exists x. \varphi x \rfloor$ 
nitpick[expect=genuine] oops — Countermodel found
lemma EmptyProperty5:  $\lfloor \varphi \neq (\lambda x::e. \perp) \rfloor \longrightarrow \lfloor \exists^E x. \varphi x \rfloor$ 
nitpick[expect=genuine] oops — Countermodel found

end

```

3.3 ModalFilter.thy (Figure 5 of [2])

Set filter and ultrafilter formalized for our modal logic setting.

```

theory ModalFilter imports HOMLinHOL
begin

type-synonym  $\tau = e \Rightarrow \sigma$ 
abbreviation Element:: $\tau \Rightarrow (\tau \Rightarrow \sigma) \Rightarrow \sigma$  (infix  $\in 90$ ) where  $\varphi \in S \equiv S \models \varphi$ 
abbreviation EmptySet:: $\tau$  ( $\emptyset$ ) where  $\emptyset \equiv \lambda x. \perp$ 
abbreviation UniversalSet:: $\tau$  ( $U$ ) where  $U \equiv \lambda x. \top$ 
abbreviation Subset:: $\tau \Rightarrow \tau \Rightarrow \sigma$  (infix  $\subseteq 80$ )
where  $\varphi \subseteq \psi \equiv \forall x. ((\varphi x) \supset (\psi x))$ 
abbreviation SubsetE:: $\tau \Rightarrow \tau \Rightarrow \sigma$  (infix  $\subseteq^E 80$ )

```

```

where  $\varphi \subseteq^E \psi \equiv \forall^E x. ((\varphi x) \supset (\psi x))$ 
abbreviation Intersection:: $\tau \Rightarrow \tau \Rightarrow \tau$  (infix  $\sqcap$  91)
  where  $\varphi \sqcap \psi \equiv \lambda x. ((\varphi x) \wedge (\psi x))$ 
abbreviation Inverse:: $\tau \Rightarrow \tau$  ( $^{-1}$ )
  where  $^{-1} \psi \equiv \lambda x. \neg(\psi x)$ 
abbreviation Filter  $\Phi \equiv \mathbf{U} \in \Phi \wedge \neg(\emptyset \in \Phi) \wedge$ 
   $(\forall \varphi \psi. \varphi \in \Phi \wedge \varphi \subseteq^E \psi \supset \psi \in \Phi) \wedge (\forall \varphi \psi. \varphi \in \Phi \wedge \psi \in \Phi \supset \varphi \sqcap \psi \in \Phi)$ 
abbreviation UFilter  $\Phi \equiv \text{Filter } \Phi \wedge (\forall \varphi. \varphi \in \Phi \vee (^{-1} \varphi) \in \Phi)$ 
abbreviation FilterP  $\Phi \equiv \mathbf{U} \in \Phi \wedge \neg(\emptyset \in \Phi) \wedge (\forall \varphi \psi. \varphi \in \Phi \wedge \varphi \subseteq \psi \supset \psi \in \Phi) \wedge$ 
   $(\forall \varphi \psi. \varphi \in \Phi \wedge \psi \in \Phi \supset \varphi \sqcap \psi \in \Phi)$ 
abbreviation UFilterP  $\Phi \equiv \text{FilterP } \Phi \wedge (\forall \varphi. \varphi \in \Phi \vee (^{-1} \varphi) \in \Phi)$ 

end

```

4 Gödels ontological argument – 1970 manuscript

4.1 GoedelVariantHOML1.thy (Figure 6 of [2])

Gödels axioms and definitions, as presented in the 1970 manuscript, are inconsistent. Actualist quantifiers (avoiding existential import) are used for quantification over entities, otherwise possibilist quantifiers are used.

```

theory GoedelVariantHOML1 imports HOMLinHOL
begin

```

```

consts PositiveProperty:: $(e \Rightarrow \sigma) \Rightarrow \sigma$  ( $P$ )
axiomatization where Ax1:  $[P \varphi \wedge P \psi \supset P (\varphi \cdot \psi)]$ 
axiomatization where Ax2a:  $[P \varphi \vee^e P \sim \varphi]$ 
definition God ( $G$ ) where  $G x \equiv \forall \varphi. P \varphi \supset \varphi x$ 
abbreviation PropertyInclusion ( $\neg \supset_N -$ ) where  $\varphi \supset_N \psi \equiv \square (\forall^E y. \varphi y \supset \psi y)$ 
definition Essence ( $\neg \text{Ess.}-$ ) where  $\varphi \text{ Ess. } x \equiv \forall \psi. \psi x \supset (\varphi \supset_N \psi)$ 
axiomatization where Ax2b:  $[P \varphi \supset \square P \varphi]$ 
lemma Ax2b':  $[\neg P \varphi \supset \square (\neg P \varphi)]$  using Ax2a Ax2b by blast
theorem Th1:  $[G x \supset G \text{ Ess. } x]$  using Ax2a Ax2b Essence-def God-def by (smt (verit))
definition NecExist ( $E$ ) where  $E x \equiv \forall \varphi. \varphi \text{ Ess. } x \supset \square (\exists^E x. \varphi x)$ 
axiomatization where Ax3:  $[P E]$ 
theorem Th2:  $[G x \supset \square (\exists^E y. G y)]$  using Ax3 Th1 God-def NecExist-def by

```

smt

theorem *Th3*: $\lfloor \diamond(\exists^E x. G x) \supset \square(\exists^E y. G y) \rfloor$

— sledgehammer(Th2 Rsymm) — Proof found

proof —

have 1: $\lfloor (\exists^E x. G x) \supset \square(\exists^E y. G y) \rfloor$ **using** Th2 **by** blast

have 2: $\lfloor \diamond(\exists^E x. G x) \supset \diamond\square(\exists^E y. G y) \rfloor$ **using** 1 **by** blast

have 3: $\lfloor \diamond(\exists^E x. G x) \supset \square(\exists^E y. G y) \rfloor$ **using** 2 Rsymm **by** blast

thus ?thesis **by** blast

qed

axiomatization where *Ax4*: $\lfloor P \varphi \wedge (\varphi \supset_N \psi) \supset P \psi \rfloor$

lemma *True nitpick*[*satisfy, expect=unknown*] **oops** — No model found

lemma *EmptyEssL*: $\lfloor (\lambda y. \perp) \text{ Ess. } x \rfloor$ **using** Essence-def **by** auto

theorem *Inconsistency: False*

— sledgehammer(Ax2a Ax3 Ax4 EmptyEssL NecExist_def) — Proof found

proof —

have 1: $\lfloor \neg(P(\lambda x. \perp)) \rfloor$ **using** Ax2a Ax4 **by** blast

have 2: $\lfloor P(\lambda x. (\lambda y. \perp)) \text{ Ess. } x \supset \square(\exists^E z. (\lambda y. \perp) z) \rfloor$ **using** Ax3 Ax4 NecExist_def **by** smt

have 3: $\lfloor P(\lambda x. \square(\exists^E z. (\lambda x. \perp) z)) \rfloor$ **using** 2 EmptyEssL **by** simp

have 4: $\lfloor P(\lambda x. \square \perp) \rfloor$ **using** 3 **by** auto

have 5: $\lfloor P(\lambda x. \perp) \rfloor$ **using** 4 Ax2a Ax4 **by** smt

have 6: $\lfloor \perp \rfloor$ **using** 1 5 **by** blast

thus ?thesis **by** blast

qed

end

4.2 GoedelVariantHOML2.thy (Figure 7 of [2])

After an appropriate modification of the definition of essence in Gödels 1970 ontological proof, the inconsistency revealed in Figure 6 is avoided, and the argument can be successfully verified in modal logic S5 (indeed, as shown, only symmetry of the accessibility relation is actually needed). Actualist quantifiers (avoiding existential import) are used for quantification over entities, otherwise possibilist quantifiers are used.

theory *GoedelVariantHOML2 imports HOMLinHOL ModalFilter*
begin

consts *PositiveProperty*:: $(e \Rightarrow \sigma) \Rightarrow \sigma$ (*P*)

axiomatization where *Ax1*: $\lfloor P \varphi \wedge P \psi \supset P (\varphi \cdot \psi) \rfloor$

axiomatization where *Ax2a*: $\lfloor P \varphi \vee^e P \sim \varphi \rfloor$

definition *God* (G) **where** $G\ x \equiv \forall \varphi. P\ \varphi \supseteq \varphi\ x$

abbreviation *PropertyInclusion* (\supseteq_N) **where** $\varphi \supseteq_N \psi \equiv \square(\forall^E y. \varphi\ y \supseteq \psi\ y)$

definition *Essence* ($\neg\text{-}\text{Ess}.$) **where** $\varphi \text{ Ess. } x \equiv \varphi\ x \wedge (\forall \psi. \psi\ x \supseteq (\varphi \supseteq_N \psi))$

axiomatization **where** *Ax2b*: $[P\ \varphi \supseteq \square\ P\ \varphi]$

lemma *Ax2b'*: $[\neg P\ \varphi \supseteq \square(\neg P\ \varphi)]$ **using** *Ax2a Ax2b* **by** *blast*

theorem *Th1*: $[G\ x \supseteq G\ \text{Ess. } x]$ **using** *Ax2a Ax2b Essence-def God-def* **by** (*smt (verit)*)

definition *NecExist* (E) **where** $E\ x \equiv \forall \varphi. \varphi \text{ Ess. } x \supseteq \square(\exists^E x. \varphi\ x)$

axiomatization **where** *Ax3*: $[P\ E]$

theorem *Th2*: $[G\ x \supseteq \square(\exists^E y. G\ y)]$ **using** *Ax3 Th1 God-def NecExist-def* **by** *smt*

theorem *Th3*: $[\diamond(\exists^E x. G\ x) \supseteq \square(\exists^E y. G\ y)]$
— sledgehammer(*Th2 Rsymm*) — Proof found

proof —
have 1: $[(\exists^E x. G\ x) \supseteq \square(\exists^E y. G\ y)]$ **using** *Th2* **by** *blast*
have 2: $[\diamond(\exists^E x. G\ x) \supseteq \diamond\square(\exists^E y. G\ y)]$ **using** 1 **by** *blast*
have 3: $[\diamond(\exists^E x. G\ x) \supseteq \square(\exists^E y. G\ y)]$ **using** 2 *Rsymm* **by** *blast*
thus ?thesis **by** *blast*
qed

axiomatization **where** *Ax4*: $[P\ \varphi \wedge (\varphi \supseteq_N \psi) \supseteq P\ \psi]$

lemma *True nitpick*[*satisfy, card=1, eval=[P (\lambda x. ⊥)]*] **oops** — One model found of cardinality one

abbreviation *PosProps* $\Phi \equiv \forall \varphi. \Phi\ \varphi \supseteq P\ \varphi$
abbreviation *ConjOfPropsFrom* $\varphi\ \Phi \equiv \square(\forall^E z. \varphi\ z \leftrightarrow (\forall \psi. \Phi\ \psi \supseteq \psi\ z))$
axiomatization **where** *Ax1Gen*: $[(\text{PosProps } \Phi \wedge \text{ConjOfPropsFrom } \varphi\ \Phi) \supseteq P\ \varphi]$

lemma *L*: $[P\ G]$ **using** *Ax1Gen God-def* **by** (*smt (verit)*)

theorem *Th4*: $[\diamond(\exists^E x. G\ x)]$ **using** *Ax2a Ax4 L* **by** *blast*

theorem *Th5*: $[\square(\exists^E x. G\ x)]$ **using** *Th3 Th4* **by** *blast*

lemma *MC*: $[\varphi \supseteq \square\varphi]$
— sledgehammer(*Ax2a Ax2b Th5 God_def Rsymm*) — Proof found

proof —

```

{fix w fix Q
  have 1:  $\forall x. (G x w \longrightarrow (\forall Z. Z x \supset \Box(\forall^E z. G z \supset Z z)) w)$  using Ax2a
Ax2b God-def by smt
  have 2:  $(\exists x. G x w) \longrightarrow ((Q \supset \Box(\forall^E z. G z \supset Q)) w)$  using 1 by force
  have 3:  $(Q \supset \Box Q) w$  using 2 Th5 Rsymm by blast}
  thus ?thesis by auto
qed

lemma PosProps:  $[P(\lambda x. \top) \wedge P(\lambda x. x = x)]$  using Ax2a Ax4 by blast
lemma NegProps:  $[\neg P(\lambda x. \perp) \wedge \neg P(\lambda x. x \neq x)]$  using Ax2a Ax4 by blast
lemma UniqueEss1:  $[\varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \Box(\forall^E y. \varphi y \leftrightarrow \psi y)]$  using
Essence-def by smt
lemma UniqueEss2:  $[\varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \Box(\varphi \equiv \psi)]$  nitpick[card i=1] oops
— Countermodel found
lemma UniqueEss3:  $[\varphi \text{ Ess. } x \supset \Box(\forall^E y. \varphi y \supset y \equiv x)]$  using Essence-def MC
by auto
lemma Monotheism:  $[G x \wedge G y \supset x \equiv y]$  using Ax2a God-def by smt
lemma Filter:  $[Filter P]$  using Ax1 Ax4 MC NegProps PosProps Rsymm by smt
lemma UltraFilter:  $[UFilter P]$  using Ax2a Filter by smt
lemma True nitpick[satisfy,card=1,eval=[P(\lambda x. \perp)]] oops — One model found
of cardinality one

end

```

4.3 GoedelVariantHOML3.thy (Figure 8 of [2])

After an appropriate modification of the notion of necessary property inclusion in Gödels 1970 ontological proof, the inconsistency revealed in Figure 6 is avoided, and the argument can be successfully verified in modal logic S5 (indeed, as shown, only symmetry of the accessibility relation is actually needed). Actualist quantifiers (avoiding existential import) are used for quantification over entities, otherwise possibilist quantifiers are used.

```

theory GoedelVariantHOML3 imports HOMLinHOL ModalFilter
begin

consts PositiveProperty:: $(e \Rightarrow \sigma) \Rightarrow \sigma (P)$ 

axiomatization where Ax1:  $[P \varphi \wedge P \psi \supset P(\varphi \cdot \psi)]$ 
axiomatization where Ax2a:  $[P \varphi \vee^e P \sim \varphi]$ 

definition God (G) where  $G x \equiv \forall \varphi. P \varphi \supset \varphi x$ 

abbreviation PropertyInclusion (- $\supset_N$ ) where  $\varphi \supset_N \psi \equiv \Box(\varphi \neq (\lambda x. \perp) \wedge$ 
 $(\forall^E y. \varphi y \supset \psi y))$ 

definition Essence (- $\text{Ess.}$ -) where  $\varphi \text{ Ess. } x \equiv \forall \psi. \psi x \supset (\varphi \supset_N \psi)$ 

```

axiomatization where $Ax2b: [P \varphi \supset \square P \varphi]$

lemma $Ax2b': [\neg P \varphi \supset \square(\neg P \varphi)]$ **using** $Ax2a$ $Ax2b$ **by** *blast*

theorem $Th1: [G x \supset G \text{Ess. } x]$ **using** $Ax2a$ $Ax2b$ *Essence-def God-def* **by** (*smt (verit)*)

definition $NecExist (E)$ **where** $E x \equiv \forall \varphi. \varphi \text{Ess. } x \supset \square(\exists^E x. \varphi x)$

axiomatization where $Ax3: [P E]$

theorem $Th2: [G x \supset \square(\exists^E y. G y)]$ **using** $Ax3$ $Th1$ *God-def NecExist-def* **by** *smt*

theorem $Th3: [\diamond(\exists^E x. G x) \supset \square(\exists^E y. G y)]$

— sledgehammer($Th2$ *Rsymm*) — Proof found

proof —

have 1: $[(\exists^E x. G x) \supset \square(\exists^E y. G y)]$ **using** $Th2$ **by** *blast*

have 2: $[\diamond(\exists^E x. G x) \supset \diamond\square(\exists^E y. G y)]$ **using** 1 **by** *blast*

have 3: $[\diamond(\exists^E x. G x) \supset \square(\exists^E y. G y)]$ **using** 2 *Rsymm* **by** *blast*

thus *?thesis by blast*

qed

axiomatization where $Ax4: [P \varphi \wedge (\varphi \supset_N \psi) \supset P \psi]$

lemma $True$ **nitpick**[*satisfy, card=1, eval=[P (\lambda x. \perp)]*] **oops** — Two models found of cardinality one

abbreviation $PosProps \Phi \equiv \forall \varphi. \Phi \varphi \supset P \varphi$

abbreviation $ConjOfPropsFrom \varphi \Phi \equiv \square(\forall^E z. \varphi z \leftrightarrow (\forall \psi. \Phi \psi \supset \psi z))$

axiomatization where $Ax1Gen: [(PosProps \Phi \wedge ConjOfPropsFrom \varphi \Phi) \supset P \varphi]$

lemma $L: [P G]$ **using** $Ax1Gen$ *God-def* **by** (*smt (verit)*)

theorem $Th4: [\diamond(\exists^E x. G x)]$

— sledgehammer[*timeout=200*]($Ax2a$ L $Ax1Gen$) **oops** — sorry — Proof found

axiomatization where $Th4: [\diamond(\exists^E x. G x)]$

theorem $Th5: [\square(\exists^E x. G x)]$ **using** $Th3$ $Th4$ **by** *blast*

lemma $MC: [\varphi \supset \square\varphi]$

— sledgehammer($Ax2a$ $Ax2b$ $Th5$ *God_def Rsymm*) — Proof found

proof — {fix w fix Q

have 1: $\forall x. (G x w \longrightarrow (\forall Z. Z x \supset \square(\forall^E z. G z \supset Z z)) w)$ **using** $Ax2a$ $Ax2b$ *God-def by smt*

have 2: $(\exists x. G x w) \longrightarrow ((Q \supset \square(\forall^E z. G z \supset Q)) w)$ **using** 1 **by** *force*

have 3: $(Q \supset \square Q) w$ **using** 2 *Th5 Rsymm* **by** *blast*}

thus *?thesis by auto*

qed

```

lemma PosProps: [ $P(\lambda x. \top) \wedge P(\lambda x. x = x)$ ] using Ax2a Ax4 L Th4 by smt
lemma NegProps: [ $\neg P(\lambda x. \perp) \wedge \neg P(\lambda x. x \neq x)$ ] using Ax2a Ax4 L Th4 by smt
lemma UniqueEss1: [ $\varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \square(\forall^E y. \varphi y \leftrightarrow \psi y)$ ] oops —
Unclear, open question
lemma UniqueEss2: [ $\varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \square(\varphi \equiv \psi)$ ] oops — Unclear, open
question
lemma UniqueEss3: [ $\varphi \text{ Ess. } x \supset \square(\forall^E y. \varphi y \supset y \equiv x)$ ] using Essence-def MC
by auto
lemma Monotheism: [ $G x \wedge G y \supset x \equiv y$ ] using Ax2a God-def by smt
lemma Filter: [ $\text{Filter } P$ ] using Ax1 Ax4 MC NegProps PosProps Rsymm by smt
lemma UltraFilter: [ $\text{UFilter } P$ ] using Ax2a Filter by smt
lemma True nitpick[satisfy,card=1,eval=[ $P(\lambda x. \top)$ ]] oops — One model found
of cardinality one

```

end

4.4 ThereIsNoEvil1.thy (Figure 10 of [2])

Importing Gödels modified axioms from Figure 7 we can prove that necessarily there exists no entity that possesses all non-positive (=negative) properties.

```

theory ThereIsNoEvil1 imports GoedelVariantHOML2
begin

```

```

definition Evil (Evil) where Evil  $x \equiv \forall \varphi. \neg P \varphi \supset \varphi x$ 

```

```

theorem NecNoEvil: [ $\square(\neg(\exists^E x. \text{ Evil } x))$ ]
— sledgehammer(Ax2a Ax4 Evil_def) — Proof found
proof —
  have [ $\neg P(\lambda y. \perp)$ ] using Ax2a Ax4 by blast
  hence [ $(\forall^E x. \text{ Evil } x \supset (\lambda y. \perp) x)$ ] using Evil-def by auto
  hence [ $(\forall^E x. \text{ Evil } x \supset \perp)$ ] by auto
  hence [ $(\exists^E x. \text{ Evil } x) \supset \perp$ ] by auto
  hence [ $\neg(\exists^E x. \text{ Evil } x)$ ] by blast
  thus ?thesis by blast
qed

```

end

4.5 ThereIsNoEvil2.thy (Figure 11 of [2])

Importing Gödels modified axioms from Figure 8 we can prove that necessarily there exists no entity that possesses all non-positive (=negative) properties.

```

theory ThereIsNoEvil2 imports GoedelVariantHOML3
begin

```

```

definition Evil (Evil) where Evil x ≡ ∀φ. ¬P φ ⊃ φ x
theorem NecNoEvil: [□(¬(∃Ex. Evil x))]
  — sledgehammer(Ax1Gen Ax2a Evil_def) oops — Proof found

end

```

5 Scotts variant

5.1 ScottVariantHOML.thy (Figure 12 of [2])

Verification of Scotts variant of Gödels ontological argument. Actualist quantifiers (avoiding existential import) are used for quantification over entities, otherwise possibilist quantifiers are used.

```

theory ScottVariantHOML imports HOMLinHOL ModalFilter
begin

consts PositiveProperty::(e⇒σ)⇒σ (P)

axiomatization where A1: [¬P φ ↔ P ∼φ]
axiomatization where A2: [P φ ∧ □(∀Ey. φ y ⊃ ψ y) ⊃ P ψ]
theorem T1: [P φ ⊃ ◇(∃Ex. φ x)] using A1 A2 by blast

definition God (G) where G x ≡ ∀φ. P φ ⊃ φ x
axiomatization where A3: [P G]
theorem Coro: [◇(∃Ex. G x)] using A3 T1 by blast
axiomatization where A4: [P φ ⊃ □ P φ]

definition Ess (-Ess.-) where φ Ess. x ≡ φ x ∧ (∀ψ. ψ x ⊃ □(∀Ey. φ y ⊃ ψ y))
theorem T2: [G x ⊃ G Ess. x] using A1 A4 Ess-def God-def by fastforce

definition NecExist (NE) where NE x ≡ ∀φ. φ Ess. x ⊃ □(∃Ex. φ x)
axiomatization where A5: [P NE]

lemma True nitpick[satisfy,card=1,eval=[P (λx.⊤)]] oops — One model found
of cardinality one

theorem T3: [□(∃Ex. G x)]
  — sledgehammer(A5 Coro God_def NecExist_def Rsymm T2) — Proof found
proof —

```

```

have 1:  $\lfloor (G x \supset NE x) \wedge (G \text{Ess. } x \supset \Box(\exists^E x. G x)) \rfloor$  using A5 Ess-def
God-def NecExist-def by smt
hence 2:  $\lfloor (\exists^E x. G x) \supset \Box(\exists^E x. G x) \rfloor$  using A5 God-def NecExist-def T2
by smt
hence 3:  $\lfloor \Diamond(\exists^E x. G x) \supset (\Diamond(\Box(\exists^E x. G x)) \supset \Box(\exists^E x. G x)) \rfloor$  using Rsymm
by blast
thus ?thesis using 2 Coro by blast
qed

lemma MC:  $\lfloor \varphi \supset \Box\varphi \rfloor$ 
— sledgehammer(A1 A4 God_def Rsymm T3) — Proof found
proof – {fix w fix Q
have 1:  $\forall x. (G x w \longrightarrow (\forall Z. Z x \supset \Box(\forall^E z. ((G z) \supset (Z z)))) w)$  using A1
A4 God-def by smt
have 2:  $(\exists x. G x w) \longrightarrow ((Q \supset \Box(\forall^E z. ((G z) \supset Q))) w)$  using 1 by force
have 3:  $(Q \supset \Box Q) w$  using 2 T3 Rsymm by blast}
thus ?thesis by auto
qed

lemma PosProps:  $\lfloor P(\lambda x. \top) \wedge P(\lambda x. x = x) \rfloor$  using A1 A2 by blast
lemma NegProps:  $\lfloor \neg P(\lambda x. \perp) \wedge \neg P(\lambda x. x \neq x) \rfloor$  using A1 A2 by blast
lemma UniqueEss1:  $\lfloor \varphi \text{Ess. } x \wedge \psi \text{Ess. } x \supset \Box(\forall^E y. \varphi y \leftrightarrow \psi y) \rfloor$  using Ess-def
by smt
lemma UniqueEss2:  $\lfloor \varphi \text{Ess. } x \wedge \psi \text{Ess. } x \supset \Box(\varphi \equiv \psi) \rfloor$  nitpick[card i=1] oops
— Countermodel found
lemma UniqueEss3:  $\lfloor \varphi \text{Ess. } x \supset \Box(\forall^E y. \varphi y \supset y \equiv x) \rfloor$  using Ess-def MC by
auto
lemma Monotheism:  $\lfloor G x \wedge G y \supset x \equiv y \rfloor$  using A1 God-def by smt
lemma Filter:  $\lfloor \text{Filter } P \rfloor$  using A1 God-def Rsymm T1 T3 by (smt (verit, best))
lemma UltraFilter:  $\lfloor \text{UFilter } P \rfloor$  using Filter A1 by blast
lemma True nitpick[satisfy,card=1,eval= $\lfloor P(\lambda x. \perp) \rfloor$ ] oops — One model found
of cardinality one

end

```

5.2 HOMLinHOLonlyK.thy (slight variation of Figure 3 of [2])

Shallow embedding of higher-order modal logic (HOML) in the classical higher-order logic (HOL) of Isabelle/HOL utilizing the LogiKEy methodology. Here logic K is introduced.

```

theory HOMLinHOLonlyK imports Main
begin

```

```

— Global parameters setting for the model finder nitpick and the parser; unimport
for the reader
nitpick-params[user-axioms,expect=genuine,show-all,format=2,max-genuine=3]
declare[[syntax-ambiguity-warning=false]]

```

— Type i is associated with possible worlds and type e with entities

typedecl i — Possible worlds

typedecl e — Individuals/entities

type-synonym $\sigma = i \Rightarrow \text{bool}$ — World-lifted propositions

type-synonym $\tau = e \Rightarrow \sigma$ — modal properties

consts $R::i \Rightarrow i \Rightarrow \text{bool}$ (-r-) — Accessibility relation between worlds

— Logical connectives (operating on truth-sets)

abbreviation $Mbot::\sigma (\perp)$ where $\perp \equiv \lambda w. \text{False}$

abbreviation $Mtop::\sigma (\top)$ where $\top \equiv \lambda w. \text{True}$

abbreviation $Mneg::\sigma \Rightarrow \sigma (\neg \cdot [52]53)$ where $\neg \varphi \equiv \lambda w. \neg(\varphi w)$

abbreviation $Mand::\sigma \Rightarrow \sigma \Rightarrow \sigma (\text{infixl} \wedge 50)$ where $\varphi \wedge \psi \equiv \lambda w. \varphi w \wedge \psi w$

abbreviation $Mor::\sigma \Rightarrow \sigma \Rightarrow \sigma (\text{infixl} \vee 49)$ where $\varphi \vee \psi \equiv \lambda w. \varphi w \vee \psi w$

abbreviation $Mimp::\sigma \Rightarrow \sigma \Rightarrow \sigma (\text{infixr} \supset 48)$ where $\varphi \supset \psi \equiv \lambda w. \varphi w \rightarrow \psi w$

abbreviation $Mequiv::\sigma \Rightarrow \sigma \Rightarrow \sigma (\text{infixl} \leftrightarrow 47)$ where $\varphi \leftrightarrow \psi \equiv \lambda w. \varphi w \leftrightarrow \psi w$

abbreviation $Mbox::\sigma \Rightarrow \sigma (\Box \cdot [54]55)$ where $\Box \varphi \equiv \lambda w. \forall v. w \mathbf{r} v \rightarrow \varphi v$

abbreviation $Mdia::\sigma \Rightarrow \sigma (\Diamond \cdot [54]55)$ where $\Diamond \varphi \equiv \lambda w. \exists v. w \mathbf{r} v \wedge \varphi v$

abbreviation $Mprimeq::'a \Rightarrow 'a \Rightarrow \sigma (=)$ where $x = y \equiv \lambda w. x = y$

abbreviation $Mprimneg::'a \Rightarrow 'a \Rightarrow \sigma (\neq)$ where $x \neq y \equiv \lambda w. x \neq y$

abbreviation $Mnegpred::\tau \Rightarrow \tau (\sim)$ where $\sim \Phi \equiv \lambda x. \lambda w. \neg \Phi x w$

abbreviation $Mconpred::\tau \Rightarrow \tau \Rightarrow \tau (\text{infixl} . 50)$ where $\Phi . \Psi \equiv \lambda x. \lambda w. \Phi x w \wedge \Psi x w$

abbreviation $Mexclor::\sigma \Rightarrow \sigma \Rightarrow \sigma (\text{infixl} \vee^e 49)$ where $\varphi \vee^e \psi \equiv (\varphi \vee \psi) \wedge \neg(\varphi \wedge \psi)$

— Possibilist quantifiers (polymorphic)

abbreviation $Mallposs::('a \Rightarrow \sigma) \Rightarrow \sigma (\forall)$ where $\forall \Phi \equiv \lambda w. \forall x. \Phi x w$

abbreviation $Mallpossb (\text{binder } \forall [8]9)$ where $\forall x. \varphi(x) \equiv \forall \varphi$

abbreviation $Mexpiposs::('a \Rightarrow \sigma) \Rightarrow \sigma (\exists)$ where $\exists \Phi \equiv \lambda w. \exists x. \Phi x w$

abbreviation $Mexpipossb (\text{binder } \exists [8]9)$ where $\exists x. \varphi(x) \equiv \exists \varphi$

— Actualist quantifiers (for individuals/entities)

consts $existsAt::e \Rightarrow \sigma (@)$

abbreviation $Mallact::(e \Rightarrow \sigma) \Rightarrow \sigma (\forall^E)$ where $\forall^E \Phi \equiv \lambda w. \forall x. x @ w \rightarrow \Phi x w$

abbreviation $Mallactb (\text{binder } \forall^E [8]9)$ where $\forall^E x. \varphi(x) \equiv \forall^E \varphi$

abbreviation $Mexist::(e \Rightarrow \sigma) \Rightarrow \sigma (\exists^E)$ where $\exists^E \Phi \equiv \lambda w. \exists x. x @ w \wedge \Phi x w$

abbreviation $Mexistb (\text{binder } \exists^E [8]9)$ where $\exists^E x. \varphi(x) \equiv \exists^E \varphi$

— Leibniz equality (polymorphic)

abbreviation $Mleibeq::'a \Rightarrow 'a \Rightarrow \sigma (=)$ where $x \equiv y \equiv \forall P. P x \supset P y$

— Meta-logical predicate for global validity

abbreviation $Mvalid::\sigma \Rightarrow \text{bool} ([\cdot])$ where $[\psi] \equiv \forall w. \psi w$

end

5.3 ScottVariantHOMLinK.thy (Figure 13 of [2])

Scotts variant of Gödels argument fails for base logic K (but only it the last step).

```
theory ScottVariantHOMLinK imports HOMLinHOLonlyK
begin

consts PositiveProperty::( $e \Rightarrow \sigma \Rightarrow \sigma$ ) ( $P$ )

axiomatization where A1: [ $\neg P \varphi \leftrightarrow P \sim \varphi$ ]

axiomatization where A2: [ $P \varphi \wedge \square(\forall^E y. \varphi y \supset \psi y) \supset P \psi$ ]

theorem T1: [ $P \varphi \supset \diamond(\exists^E x. \varphi x)$ ] using A1 A2 by blast

definition God (G) where G x ≡  $\forall \varphi. P \varphi \supset \varphi x$ 

axiomatization where A3: [ $P G$ ]

theorem Coro: [ $\diamond(\exists^E x. G x)$ ] nitpick[satisfy,eval=G] using A3 T1 by blast

axiomatization where A4: [ $P \varphi \supset \square P \varphi$ ]

definition Ess (-Ess.-) where φ Ess. x ≡  $\varphi x \wedge (\forall \psi. \psi x \supset \square(\forall^E y. \varphi y \supset \psi y))$ 

theorem T2: [ $G x \supset G \text{Ess. } x$ ] using A1 A4 Ess-def God-def by fastforce

definition NecExist (NE) where NE x ≡  $\forall \varphi. \varphi \text{Ess. } x \supset \square(\exists^E x. \varphi x)$ 

axiomatization where A5: [ $P NE$ ]

lemma True nitpick[satisfy,card=1,eval=[ $P (\lambda x. \top)$ ]] oops — One model found of cardinality one

theorem T3: [ $\square(\exists^E x. G x)$ ] nitpick[card e=1, card i=2, eval=G] oops — Counterexample

lemma MC: [ $\varphi \supset \square \varphi$ ] nitpick[card e=1, card i=2, eval=G] oops — Counterexample

end
```

5.4 ScottVariantHOMLAndersonQuant.thy (Figure 15 of [2])

Verification of Scotts variant of Gödels argument with a mixed use of actutilist and possibilist quantifiers for entities; cf. Footnote 20 in [2].

```
theory ScottVariantHOMLAndersonQuant imports HOMLinHOL ModalFilter
```

```

begin

consts PositiveProperty::( $e \Rightarrow \sigma \Rightarrow \sigma$ ) ( $P$ )

axiomatization where A1: [ $\neg P \varphi \leftrightarrow P \sim \varphi$ ]

axiomatization where A2: [ $P \varphi \wedge \square(\forall y. \varphi y \supset \psi y) \supset P \psi$ ]

theorem T1: [ $P \varphi \supset \diamond(\exists x. \varphi x)$ ] using A1 A2 by smt

definition God (G) where G x ≡  $\forall \varphi. P \varphi \supset \varphi x$ 

axiomatization where A3: [ $P G$ ]

theorem Coro: [ $\diamond(\exists x. G x)$ ] using A3 T1 by blast

axiomatization where A4: [ $P \varphi \supset \square P \varphi$ ]

definition Ess (-Ess-) where φ Ess. x ≡  $\varphi x \wedge (\forall \psi. \psi x \supset \square(\forall y::e. \varphi y \supset \psi y))$ 

theorem T2: [ $G x \supset G \text{Ess. } x$ ] using A1 A4 Ess-def God-def by smt

definition NecExist (NE) where NE x ≡  $\forall \varphi. \varphi \text{Ess. } x \supset \square(\exists^E x. \varphi x)$ 

axiomatization where A5: [ $P \text{NE}$ ]

lemma True nitpick[satisfy,card=1,eval=] $[P (\lambda x. \top)]$  oops — One model found of cardinality one

theorem T3: [ $\square(\exists^E x. G x)$ ]
— sledgehammer(A5 Coro God_def NecExist_def Rsymm T2) — Proof found
proof —
  have 1: [ $(G x \supset \text{NE } x) \wedge (G \text{Ess. } x \supset \square(\exists^E x. G x))$ ] using A5 Ess-def God-def NecExist-def by smt
  hence 2: [ $(\exists x. G x) \supset \square(\exists^E x. G x)$ ] using A5 God-def NecExist-def T2 by smt
  hence 3: [ $\diamond(\exists x. G x) \supset (\diamond(\square(\exists x. G x)) \supset \square(\exists^E x. G x))$ ] using Rsymm by blast
  thus ?thesis using 2 Coro by blast
qed

lemma MC: [ $\varphi \supset \square \varphi$ ]
— sledgehammer(A1 A4 God_def Rsymm T3) — Proof found
proof — {fix w fix Q
  have 1:  $\forall x. (G x w \longrightarrow (\forall Z. Z x \supset \square(\forall z. ((G z) \supset (Z z)))) w)$  using A1 A4 God-def by smt
  have 2:  $(\exists x. G x w) \longrightarrow ((Q \supset \square(\forall z. ((G z) \supset Q))) w)$  using 1 by force
  have 3:  $(Q \supset \square Q) w$  using 2 T3 Rsymm by blast}

```

```

thus ?thesis by auto
qed

lemma PosProps: [ $P(\lambda x. \top) \wedge P(\lambda x. x = x)$ ] using A1 A2 by blast
lemma NegProps: [ $\neg P(\lambda x. \perp) \wedge \neg P(\lambda x. x \neq x)$ ] using A1 A2 by blast
lemma UniqueEss1: [ $\varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \square(\forall y. \varphi y \leftrightarrow \psi y)$ ] using Ess-def by smt
lemma UniqueEss2: [ $\varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \square(\varphi = \psi)$ ] nitpick[card i=2] oops — Countermodel found
lemma UniqueEss3: [ $\varphi \text{ Ess. } x \supset \square(\forall y. \varphi y \supset y \equiv x)$ ] using Ess-def MC by auto
lemma Monotheism: [ $G x \wedge G y \supset x \equiv y$ ] using A1 God-def by smt
lemma Filter: [ $\text{Filter } P$ ] using A1 God-def Rsymm T1 T3 by (smt (verit, best))
lemma UltraFilter: [ $\text{UFilter } P$ ] using Filter A1 by blast
lemma True nitpick[satisfy,card=1,eval=[ $P(\lambda x. \perp)$ ]] oops — One model found of cardinality one
end

```

6 Appendix

6.1 GoedelVariantHOML1poss.thy (Figure 16 of [2])

Gödels axioms and definitions, as presented in the 1970 manuscript, are inconsistent. In contrast to Figure 6 we here use only possibilist quantifiers and still derive falsity.

```

theory GoedelVariantHOML1poss imports HOMLinHOL
begin

consts PositiveProperty::( $e \Rightarrow \sigma \Rightarrow \sigma$ ) ( $P$ )

axiomatization where Ax1: [ $P \varphi \wedge P \psi \supset P(\varphi \cdot \psi)$ ]

axiomatization where Ax2a: [ $P \varphi \vee^e P \sim \varphi$ ]

definition God (G) where  $G x \equiv \forall \varphi. P \varphi \supset \varphi x$ 

abbreviation PropertyInclusion ( $\neg\supset_N$ ) where  $\varphi \supset_N \psi \equiv \square(\forall y::e. \varphi y \supset \psi y)$ 

definition Essence (-Ess.-) where  $\varphi \text{ Ess. } x \equiv \forall \psi. \psi x \supset (\varphi \supset_N \psi)$ 

axiomatization where Ax2b: [ $P \varphi \supset \square P \varphi$ ]

lemma Ax2b': [ $\neg P \varphi \supset \square(\neg P \varphi)$ ] using Ax2a Ax2b by blast

theorem Th1: [ $G x \supset G \text{ Ess. } x$ ] using Ax2a Ax2b Essence-def God-def by (smt (verit))

```

```

definition NecExist (E) where E x ≡ ∀φ. φ Ess. x ⊃ □(∃x. φ x)

axiomatization where Ax3: [P E]

theorem Th2: [G x ⊃ □(∃y. G y)] using Ax3 Th1 God-def NecExist-def by smt

theorem Th3: [◇(∃x. G x) ⊃ □(∃y. G y)]
— sledgehammer(Th2 Rsymm) — Proof found
proof —
  have 1: [□(∃x. G x) ⊃ □(∃y. G y)] using Th2 by blast
  have 2: [◇(∃x. G x) ⊃ ◇□(∃y. G y)] using 1 by blast
  have 3: [◇(∃x. G x) ⊃ □(∃y. G y)] using 2 Rsymm by blast
  thus ?thesis by blast
qed

axiomatization where Ax4: [P φ ∧ (φ ⊃_N ψ) ⊃ P ψ]

lemma True nitpick[satisfy,expect=unknown] oops — No model found

lemma EmptyEssL: [(\λy.⊥) Ess. x] using Essence-def by metis

theorem Inconsistency: False
— sledgehammer(Ax2a Ax3 Ax4 EmptyEssL NecExist_def) — Proof found
proof —
  have 1: [¬(P (λx.⊥))] using Ax2a Ax4 by blast
  have 2: [P (λx.(λy.⊥) Ess. x ⊃ □(∃z::e.(λy.⊥)z))] using Ax3 Ax4 NecExist-def by smt
  have 3: [P (λx.□(∃z. (λx.⊥) z))] using 2 EmptyEssL Ax4 by smt
  have 4: [P (λx.□⊥)] using 3 by auto
  have 5: [P (λx.⊥)] using 4 Ax2a Ax4 by smt
  have 6: [⊥] using 1 5 by blast
  thus ?thesis by blast
qed

end

```

6.2 GoedelVariantHOML2poss.thy (Figure 17 of [2])

After an appropriate modification of the definition of essence, the inconsistency revealed in Figure 16 is avoided, and the argument can be successfully verified in modal logic S5 (indeed, as shown, only the modal schema B is actually needed). In contrast to Figure 7 we here use only possibilist quantifiers to obtain these results.

```

theory GoedelVariantHOML2poss imports HOMLinHOL ModalFilter
begin

```

```

consts PositiveProperty::(e⇒σ)⇒σ (P)

```

axiomatization where $Ax1: [P \varphi \wedge P \psi \supset P(\varphi \cdot \psi)]$

axiomatization where $Ax2a: [P \varphi \vee^e P \sim \varphi]$

definition $God(G)$ **where** $G x \equiv \forall \varphi. P \varphi \supset \varphi x$

abbreviation $PropertyInclusion(-\supset_N-)$ **where** $\varphi \supset_N \psi \equiv \square(\forall y::e. \varphi y \supset \psi y)$

definition $Essence(-Ess.-)$ **where** $\varphi Ess. x \equiv \varphi x \wedge (\forall \psi. \psi x \supset (\varphi \supset_N \psi))$

axiomatization where $Ax2b: [P \varphi \supset \square P \varphi]$

lemma $Ax2b': [\neg P \varphi \supset \square(\neg P \varphi)]$ **using** $Ax2a$ $Ax2b$ **by** *blast*

theorem $Th1: [G x \supset G Ess. x]$ **using** $Ax2a$ $Ax2b$ *Essence-def* *God-def* **by** (*smt* (*verit*))

definition $NecExist(E)$ **where** $E x \equiv \forall \varphi. \varphi Ess. x \supset \square(\exists x. \varphi x)$

axiomatization where $Ax3: [P E]$

theorem $Th2: [G x \supset \square(\exists y. G y)]$ **using** $Ax3$ $Th1$ *God-def* *NecExist-def* **by** *smt*

theorem $Th3: [\diamond(\exists x. G x) \supset \square(\exists y. G y)]$
— sledgehammer(*Th2 Rsymm*) — Proof found

proof —
have 1: $[(\exists x. G x) \supset \square(\exists y. G y)]$ **using** *Th2* **by** *blast*
have 2: $[\diamond(\exists x. G x) \supset \diamond\square(\exists y. G y)]$ **using** 1 **by** *blast*
have 3: $[\diamond(\exists x. G x) \supset \square(\exists y. G y)]$ **using** 2 *Rsymm* **by** *blast*
thus ?*thesis* **by** *blast*

qed

axiomatization where $Ax4: [P \varphi \wedge (\varphi \supset_N \psi) \supset P \psi]$

lemma $True$ **nitpick**[*satisfy, card=1, eval=[P (\lambda x. \perp)]*] **oops** — One model found of cardinality one

abbreviation $PosProps \Phi \equiv \forall \varphi. \Phi \varphi \supset P \varphi$
abbreviation $ConjOfPropsFrom \varphi \Phi \equiv \square(\forall z. \varphi z \leftrightarrow (\forall \psi. \Phi \psi \supset \psi z))$
axiomatization where $Ax1Gen: [(PosProps \Phi \wedge ConjOfPropsFrom \varphi \Phi) \supset P \varphi]$

lemma $L: [P G]$ **using** $Ax1Gen$ *God-def* **by** *smt*

theorem $Th4: [\diamond(\exists x. G x)]$ **using** $Ax2a$ $Ax4$ L **by** *blast*

theorem $Th5: [\square(\exists x. G x)]$ **using** $Th3$ $Th4$ **by** *blast*

lemma $MC: [\varphi \supset \square \varphi]$

```

— sledgehammer(Ax2a Ax2b Th5 God_def Rsymm) — Proof found
proof — {fix w fix Q
  have 1:  $\forall x.(G x w \longrightarrow (\forall Z. Z x \supset \Box(\forall z. G z \supset Z z)) w)$  using Ax2a Ax2b
  God-def by smt
  have 2:  $(\exists x. G x w) \longrightarrow ((Q \supset \Box(\forall z. G z \supset Q)) w)$  using 1 by force
  have 3:  $(Q \supset \Box(Q)) w$  using 2 Th5 Rsymm by blast}
  thus ?thesis by auto
qed

lemma PosProps: [ $P(\lambda x. \top) \wedge P(\lambda x. x = x)$ ] using Ax2a Ax4 by blast
lemma NegProps: [ $\neg P(\lambda x. \perp) \wedge \neg P(\lambda x. x \neq x)$ ] using Ax2a Ax4 by blast
lemma UniqueEss1: [ $\varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \Box(\forall y. \varphi y \leftrightarrow \psi y)$ ] using Essence-def
  by smt
lemma UniqueEss2: [ $\varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \Box(\varphi \equiv \psi)$ ] nitpick[card i=2] oops
  — Countermodel found
lemma UniqueEss3: [ $\varphi \text{ Ess. } x \supset \Box(\forall y. \varphi y \supset y \equiv x)$ ] using Essence-def MC
  by auto
lemma Monotheism: [ $G x \wedge G y \supset x \equiv y$ ] using Ax2a God-def by smt
lemma Filter: [ $\text{FilterP } P$ ] using Ax1 Ax4 MC NegProps PosProps Rsymm by smt
lemma UltraFilter: [ $\text{UFilterP } P$ ] using Ax2a Filter by smt
lemma True nitpick[satisfy,card=1,eval=[P(\lambda x. \perp)]] oops — One model found
  of cardinality one

end

```

6.3 GoedelVariantHOML3poss.thy (Figure 18 of [2])

After an appropriate modification of the definition of necessary property implication, the inconsistency shown in Figure 16 is avoided, and the argument can be successfully verified. As shown here, this still holds when using possibilist quantifiers only.

```
theory GoedelVariantHOML3poss imports HOMLinHOL ModalFilter
begin
```

```
consts PositiveProperty::( $e \Rightarrow \sigma \Rightarrow \sigma(P)$ )
```

```
axiomatization where Ax1: [ $P \varphi \wedge P \psi \supset P(\varphi \cdot \psi)$ ]
```

```
axiomatization where Ax2a: [ $P \varphi \vee^e P \sim \varphi$ ]
```

```
definition God (G) where  $G x \equiv \forall \varphi. P \varphi \supset \varphi x$ 
```

```
abbreviation PropertyInclusion (- $\supset_N$ -) where  $\varphi \supset_N \psi \equiv \Box(\varphi \neq (\lambda x. \perp) \wedge (\forall y::e. \varphi y \supset \psi y))$ 
```

```
definition Essence (- $\text{Ess.}$ -) where  $\varphi \text{ Ess. } x \equiv \forall \psi. \psi x \supset (\varphi \supset_N \psi)$ 
```

axiomatization where $Ax2b: [P \varphi \supset \square P \varphi]$

lemma $Ax2b': [\neg P \varphi \supset \square(\neg P \varphi)]$ **using** $Ax2a$ $Ax2b$ **by** *blast*

theorem $Th1: [G x \supset G \text{Ess. } x]$ **using** $Ax2a$ $Ax2b$ *Essence-def God-def* **by** (*smt (verit)*)

definition $NecExist (E)$ **where** $E x \equiv \forall \varphi. (\varphi \text{Ess. } x) \supset \square(\exists x. \varphi x)$

axiomatization where $Ax3: [P E]$

theorem $Th2: [G x \supset \square(\exists y. G y)]$ **using** $Ax3$ $Th1$ *God-def NecExist-def* **by** *smt*

theorem $Th3: [\diamond(\exists x. G x) \supset \square(\exists y. G y)]$

— sledgehammer($Th2$ *Rsymm*) — Proof found

proof —

have 1: $[(\exists x. G x) \supset \square(\exists y. G y)]$ **using** $Th2$ **by** *blast*

have 2: $[\diamond(\exists x. G x) \supset \diamond\square(\exists y. G y)]$ **using** 1 **by** *blast*

have 3: $[\diamond(\exists x. G x) \supset \square(\exists y. G y)]$ **using** 2 *Rsymm* **by** *blast*

thus *?thesis by blast*

qed

axiomatization where $Ax4: [P \varphi \wedge (\varphi \supset_N \psi) \supset P \psi]$

lemma $True$ **nitpick**[*satisfy, card=1, eval=[P (\lambda x. \perp)]*] **oops** — One model found of cardinality one

abbreviation $PosProps \Phi \equiv \forall \varphi. \Phi \varphi \supset P \varphi$

abbreviation $ConjOfPropsFrom \varphi \Phi \equiv \square(\forall z. \varphi z \leftrightarrow (\forall \psi. \Phi \psi \supset \psi z))$

axiomatization where $Ax1Gen: [(PosProps \Phi \wedge ConjOfPropsFrom \varphi \Phi) \supset P \varphi]$

lemma $L: [P G]$ **using** $Ax1Gen$ *God-def* **by** (*smt (verit)*)

theorem $Th4: [\diamond(\exists x. G x)]$

— sledgehammer[*timeout=200*]($Ax2a$ L $Ax1Gen$) **oops** — sorry — Proof found

axiomatization where $Th4: [\diamond(\exists x. G x)]$

theorem $Th5: [\square(\exists x. G x)]$ **using** $Th3$ $Th4$ **by** *blast*

lemma $MC: [\varphi \supset \square\varphi]$

— sledgehammer($Ax2a$ $Ax2b$ $Th5$ *God_def Rsymm*) — Proof found

proof — {fix w fix Q

have 1: $\forall x. (G x w \longrightarrow (\forall Z. Z x \supset \square(\forall z. G z \supset Z z)) w)$ **using** $Ax2a$ $Ax2b$ *God-def by smt*

have 2: $(\exists x. G x w) \longrightarrow ((Q \supset \square(\forall z. G z \supset Q)) w)$ **using** 1 **by** *force*

have 3: $(Q \supset \square Q)$ **using** 2 $Th5$ *Rsymm* **by** *blast*}

thus *?thesis by auto*

qed

lemma *PosProps*: $[P(\lambda x. \top) \wedge P(\lambda x. x = x)]$ **using** *Ax2a Ax4 L Th4* **by** *smt*
lemma *NegProps*: $[\neg P(\lambda x. \perp) \wedge \neg P(\lambda x. x \neq x)]$ **using** *Ax2a Ax4 L Th4* **by** *smt*
lemma *UniqueEss1*: $[\varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \square(\forall y. \varphi y \leftrightarrow \psi y)]$ **oops** — Unclear, open question
lemma *UniqueEss2*: $[\varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \square(\varphi \equiv \psi)]$ **oops** — Unclear, open question
lemma *UniqueEss3*: $[\varphi \text{ Ess. } x \supset \square(\forall y. \varphi y \supset y \equiv x)]$ **using** *Essence-def MC by auto*
lemma *Monotheism*: $[G x \wedge G y \supset x \equiv y]$ **using** *Ax2a God-def* **by** *smt*
lemma *Filter*: $[FilterP P]$ **using** *Ax1 Ax4 MC NegProps PosProps Rsymm* **by** *smt*
lemma *UltraFilter*: $[UFilterP P]$ **using** *Ax2a Filter* **by** *smt*
lemma *True nitpick*[*satisfy, card=1, eval=[P(\lambda x. \top)]*] **oops** — One model found of cardinality one

end

6.4 ScottVariantHOMLposs.thy (Figure 19 of [2])

Scotts variant of Gödels ontological proof is still valid when using possibilist quantifiers only.

theory *ScottVariantHOMLposs* **imports** *HOMLinHOL ModalFilter*
begin

consts *PositiveProperty*:: $(e \Rightarrow \sigma) \Rightarrow \sigma (P)$
axiomatization where *A1*: $[\neg P \varphi \leftrightarrow P \sim \varphi]$
axiomatization where *A2*: $[P \varphi \wedge \square(\forall y. \varphi y \supset \psi y) \supset P \psi]$
theorem *T1*: $[P \varphi \supset \diamond(\exists x. \varphi x)]$ **using** *A1 A2* **by** *blast*
definition *God* (*G*) **where** $G x \equiv \forall \varphi. P \varphi \supset \varphi x$
axiomatization where *A3*: $[P G]$
theorem *Coro*: $[\diamond(\exists x. G x)]$ **using** *A3 T1* **by** *blast*
axiomatization where *A4*: $[P \varphi \supset \square P \varphi]$
definition *Ess (-Ess-)* **where** $\varphi \text{ Ess. } x \equiv \varphi x \wedge (\forall \psi. \psi x \supset \square(\forall y::e. \varphi y \supset \psi y))$
theorem *T2*: $[G x \supset G \text{ Ess. } x]$ **using** *A1 A4 Ess-def God-def* **by** *fastforce*
definition *NecExist* (*NE*) **where** $NE x \equiv \forall \varphi. \varphi \text{ Ess. } x \supset \square(\exists x. \varphi x)$

axiomatization where A5: $\lfloor P \text{ NE} \rfloor$

lemma $\text{True nitpick[satisfy,card=1,eval=}\lfloor P (\lambda x. \perp) \rfloor\text{]} \text{ oops}$ — One model found of cardinality one

theorem T3: $\lfloor \square(\exists x. G x) \rfloor$

— sledgehammer(A5 Coro God_def NecExist_def Rsymm T2) — Proof found
proof —

have 1: $\lfloor (G x \supset \text{NE } x) \wedge (G \text{ Ess. } x \supset \square(\exists x. G x)) \rfloor$ **using** A5 Ess-def God-def NecExist-def **by** smt

hence 2: $\lfloor (\exists x. G x) \supset \square(\exists x. G x) \rfloor$ **using** A5 God-def NecExist-def T2 **by** smt

hence 3: $\lfloor \diamond(\exists x. G x) \supset (\diamond(\square(\exists x. G x)) \supset \square(\exists x. G x)) \rfloor$ **using** Rsymm **by** blast

thus ?thesis **using** 2 Coro **by** blast

qed

lemma MC: $\lfloor \varphi \supset \square\varphi \rfloor$

— sledgehammer(A1 A4 God_def Rsymm T3) — Proof found

proof — {fix w fix Q

have 1: $\forall x. (G x w \longrightarrow (\forall Z. Z x \supset \square(\forall z. ((G z) \supset (Z z)))) w)$ **using** A1 A4 God-def **by** smt

have 2: $\lfloor (\exists x. G x w) \longrightarrow ((Q \supset \square(\forall z. ((G z) \supset Q))) w) \rfloor$ **using** 1 **by** force

have 3: $\lfloor (Q \supset \square Q) w \rfloor$ **using** 2 T3 Rsymm **by** blast}

thus ?thesis **by** auto

qed

lemma PosProps: $\lfloor P (\lambda x. \top) \wedge P (\lambda x. x = x) \rfloor$ **using** A1 A2 **by** blast

lemma NegProps: $\lfloor \neg P(\lambda x. \perp) \wedge \neg P(\lambda x. x \neq x) \rfloor$ **using** A1 A2 **by** blast

lemma UniqueEss1: $\lfloor \varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \square(\forall y. \varphi y \leftrightarrow \psi y) \rfloor$ **using** Ess-def **by** smt

lemma UniqueEss2: $\lfloor \varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \square(\varphi \equiv \psi) \rfloor$ **nitpick[card i=2] oops**
 — Countermodel found

lemma UniqueEss3: $\lfloor \varphi \text{ Ess. } x \supset \square(\forall y. \varphi y \supset y \equiv x) \rfloor$ **using** Ess-def MC **by** auto

lemma Monotheism: $\lfloor G x \wedge G y \supset x \equiv y \rfloor$ **using** A1 God-def **by** smt

lemma Filter: $\lfloor \text{FilterP } P \rfloor$ **using** A1 God-def Rsymm T1 T3 **by** (smt (verit, best))

lemma UltraFilter: $\lfloor \text{UFilterP } P \rfloor$ **using** Filter A1 **by** blast

lemma True nitpick[satisfy,card=1,eval=] $\lfloor P (\lambda x. \perp) \rfloor$ **oops** — One model found of cardinality one

end

6.5 EvilDerivable.thy (Figure 20 of [2])

The necessary existence of an Evil-like entity proved from (controversially) modified assumptions. By rejecting Gödels assumptions and instead postulating corresponding negative versions of them, as shown in the Figure 20, the necessary existence of Evil becomes derivable. The non-positive prop-

erties of this Evil-like entity are however identical to the positive properties of Gödels God-like entity.

```

theory EvilDerivable imports HOMLinHOL ModalFilter
begin

consts PositiveProperty::( $e \Rightarrow \sigma \Rightarrow \sigma$ ) $(P)$ 

definition Evil ( $Evil$ ) where  $Evil\ x \equiv \forall \varphi. \neg P\ \varphi \supset \varphi\ x$ 

definition Essence (-Ess.-) where  $\varphi\ Ess.\ x \equiv \varphi\ x \wedge (\forall \psi. \psi\ x \supset \square(\forall^E y. \varphi\ y \supset \psi\ y))$ 

definition NecExist ( $E$ ) where  $E\ x \equiv \forall \varphi. \varphi\ Ess.\ x \supset \square(\exists^E x. \varphi\ x)$ 

axiomatization where A1:  $[\neg P\ \varphi \leftrightarrow P\ \sim\varphi]$ 

axiomatization where A2:  $[\neg P\ \varphi \wedge \square(\forall^E y. \varphi\ y \supset \psi\ y) \supset \neg P\ \psi]$ 

axiomatization where A4:  $[\neg P\ Evil]$ 

axiomatization where A3:  $[\neg P\ \varphi \supset \square(\neg P\ \varphi)]$ 

axiomatization where A5:  $[\neg P\ E]$ 

lemma True nitpick[satisfy,card i=1,eval=[ $P(\lambda x.\perp)$ ],eval=[ $P(\lambda x.\top)$ ]] oops
— Model found

theorem T1:  $[\neg P\ \varphi \supset \diamond(\exists^E x. \varphi\ x)]$  using A1 A2 by blast

theorem T2:  $[\diamond(\exists^E x. Evil\ x)]$  using A4 T1 by blast

theorem T3:  $[Evil\ x \supset Evil\ Ess.\ x]$  using A1 A3 Essence-def Evil-def by (smt (verit, best))

theorem T4:  $[\diamond(\exists^E x. Evil\ x) \supset \square(\exists^E y. Evil\ y)]$  using A5 Evil-def NecExist-def Rsymm T3 by smt

theorem T5:  $[\square(\exists^E x. Evil\ x)]$  using T2 T4 by presburger

lemma MC:  $[\varphi \supset \square\varphi]$ 
— sledgehammer(A1 A3 T5 Evil_def Rsymm) oops — proof found

lemma PosProps:  $[P(\lambda x.\perp) \wedge P(\lambda x. x \neq x)]$  using A1 A2 by blast
lemma NegProps:  $[\neg P(\lambda x.\top) \wedge \neg P(\lambda x. x = x)]$  using A1 A2 by blast
lemma UniqueEss1:  $[\varphi\ Ess.\ x \wedge \psi\ Ess.\ x \supset \square(\forall^E y. \varphi\ y \leftrightarrow \psi\ y)]$  using Essence-def by smt
lemma UniqueEss2:  $[\varphi\ Ess.\ x \wedge \psi\ Ess.\ x \supset \square(\varphi = \psi)]$  nitpick[card i=2] oops
— Countermodel found
lemma Monoevilism:  $[Evil\ x \wedge Evil\ y \supset x \equiv y]$  using A1 Evil-def by smt

```

```

lemma Filter: [ $\text{Filter}(\lambda \varphi. \neg P \varphi)$ ] using A1 Evil-def Rsymm T1 T5 by (smt (verit, best))
lemma UltraFilter: [ $\text{UltraFilter}(\lambda \varphi. \neg P \varphi)$ ] using Filter A1 by blast
end

```

7 Further Appendices

7.1 GoedelVariantHOML1AndersonQuan.thy

The same as GoedelVariantHOML1, but now for a mixed use of actualist and possibilist quantifiers for entities; cf. Footnote 20 in [2].

```

theory GoedelVariantHOML1AndersonQuant imports HOMLinHOL
begin

```

```

consts PositiveProperty::( $e \Rightarrow \sigma \Rightarrow \sigma (P)$ )

```

```

axiomatization where Ax1: [ $P \varphi \wedge P \psi \supset P (\varphi \cdot \psi)$ ]

```

```

axiomatization where Ax2a: [ $P \varphi \vee^e P \sim \varphi$ ]

```

```

definition God (G) where G x  $\equiv \forall \varphi. P \varphi \supset \varphi x$ 

```

```

abbreviation PropertyInclusion ( $\neg \supset_N -$ ) where  $\varphi \supset_N \psi \equiv \square(\forall y::e. \varphi y \supset \psi y)$ 

```

```

definition Essence (-Ess.-) where  $\varphi \text{ Ess. } x \equiv \forall \psi. \psi x \supset (\varphi \supset_N \psi)$ 

```

```

axiomatization where Ax2b: [ $P \varphi \supset \square P \varphi$ ]

```

```

lemma Ax2b': [ $\neg P \varphi \supset \square(\neg P \varphi)$ ] using Ax2a Ax2b by blast

```

```

theorem Th1: [ $G x \supset G \text{ Ess. } x$ ] using Ax2a Ax2b Essence-def God-def by (smt (verit))

```

```

definition NecExist (E) where E x  $\equiv \forall \varphi. \varphi \text{ Ess. } x \supset \square(\exists^E x. \varphi x)$ 

```

```

axiomatization where Ax3: [ $P E$ ]

```

```

theorem Th2: [ $G x \supset \square(\exists^E y. G y)$ ] using Ax3 Th1 God-def NecExist-def by smt

```

```

theorem Th3: [ $\diamond(\exists^E x. G x) \supset \square(\exists^E y. G y)$ ]

```

— sledgehammer(Th2 Rsymm) — Proof found

proof —

have 1: [$(\exists^E x. G x) \supset \square(\exists^E y. G y)$] **using** Th2 **by** blast

have 2: [$\diamond(\exists^E x. G x) \supset \diamond\square(\exists^E y. G y)$] **using** 1 **by** blast

have 3: [$\diamond(\exists^E x. G x) \supset \square(\exists^E y. G y)$] **using** 2 Rsymm **by** blast

thus ?thesis **by** blast

qed

axiomatization where $Ax4: [P \varphi \wedge (\varphi \supset_N \psi) \supset P \psi]$

lemma *True* **nitpick**[*satisfy,expect=unknown*] **oops** — No model found

lemma *EmptyEssL*: $[(\lambda y. \perp) \text{ Ess. } x]$ **using** *Essence-def* **by** *auto*

theorem *Inconsistency: False*

— sledgehammer(*Ax2a Ax3 Ax4 EmptyEssL NecExist_def*) — Proof found

proof —

have 1: $[\neg(P(\lambda x. \perp))]$ **using** *Ax2a Ax4* **by** *blast*

have 2: $[P(\lambda x. (\lambda y. \perp) \text{ Ess. } x \supset \square(\exists^E z. (\lambda y. \perp) z))]$ **using** *Ax3 Ax4 NecExist-def by (smt (verit))*

have 3: $[P(\lambda x. \square(\exists^E z. (\lambda x. \perp) z))]$ **using** 2 *EmptyEssL* **by** *simp*

have 4: $[P(\lambda x. \square \perp)]$ **using** 3 **by** *auto*

have 5: $[P(\lambda x. \perp)]$ **using** 4 *Ax2a Ax4* **by** *smt*

have 6: $[\perp]$ **using** 1 5 **by** *blast*

thus *?thesis* **by** *blast*

qed

end

7.2 GoedelVariantHOML2AndersonQuan.thy

The same as GoedelVariantHOML2, but now for a mixed use of actualist and possibilist quantifiers for entities; cf. Footnote 20 in [2].

theory *GoedelVariantHOML2AndersonQuant* **imports** *HOMLinHOL ModalFilter*
begin

consts *PositiveProperty*:: $(e \Rightarrow \sigma) \Rightarrow \sigma$ (P)

axiomatization where $Ax1: [P \varphi \wedge P \psi \supset P(\varphi \cdot \psi)]$

axiomatization where $Ax2a: [P \varphi \vee^e P \sim \varphi]$

definition *God* (G) **where** $G x \equiv \forall \varphi. P \varphi \supset \varphi x$

abbreviation *PropertyInclusion* ($\neg \supset_N -$) **where** $\varphi \supset_N \psi \equiv \square(\forall y::e. \varphi y \supset \psi y)$

definition *Essence* ($\neg \text{Ess.} -$) **where** $\varphi \text{ Ess. } x \equiv \varphi x \wedge (\forall \psi. \psi x \supset (\varphi \supset_N \psi))$

axiomatization where $Ax2b: [P \varphi \supset \square P \varphi]$

lemma $Ax2b': [\neg P \varphi \supset \square(\neg P \varphi)]$ **using** *Ax2a Ax2b* **by** *blast*

theorem *Th1*: $[G x \supset G \text{ Ess. } x]$ **using** *Ax2a Ax2b Essence-def God-def by (smt (verit))*

definition *NecExist* (*E*) **where** $E\ x \equiv \forall \varphi. \varphi\ Ess.\ x \supset \square(\exists^E x. \varphi\ x)$

axiomatization where *Ax3*: $\lfloor P\ E \rfloor$

theorem *Th2*: $\lfloor G\ x \supset \square(\exists^E y. G\ y) \rfloor$ **using** *Ax3 Th1 God-def NecExist-def* **by** *smt*

theorem *Th3*: $\lfloor \diamond(\exists^E x. G\ x) \supset \square(\exists^E y. G\ y) \rfloor$

— sledgehammer(*Th2 Rsymmm*) — Proof found

proof —

have 1: $\lfloor (\exists^E x. G\ x) \supset \square(\exists^E y. G\ y) \rfloor$ **using** *Th2 by blast*

have 2: $\lfloor \diamond(\exists^E x. G\ x) \supset \diamond\square(\exists^E y. G\ y) \rfloor$ **using** 1 **by** *blast*

have 3: $\lfloor \diamond(\exists^E x. G\ x) \supset \square(\exists^E y. G\ y) \rfloor$ **using** 2 *Rsymm by blast*

thus *?thesis by blast*

qed

axiomatization where *Ax4*: $\lfloor P\ \varphi \wedge (\varphi \supset_N \psi) \supset P\ \psi \rfloor$

lemma *True nitpick*[*satisfy, card=1, eval=[P (\lambda x. \perp)]*] **oops** — One model found of cardinality one

abbreviation *PosProps* $\Phi \equiv \forall \varphi. \Phi\ \varphi \supset P\ \varphi$

abbreviation *ConjOfPropsFrom* $\varphi\ \Phi \equiv \square(\forall z. \varphi\ z \leftrightarrow (\forall \psi. \Phi\ \psi \supset \psi\ z))$

axiomatization where *Ax1Gen*: $\lfloor (\text{PosProps}\ \Phi \wedge \text{ConjOfPropsFrom}\ \varphi\ \Phi) \supset P\ \varphi \rfloor$

lemma *L*: $\lfloor P\ G \rfloor$ **using** *Ax1Gen God-def* **by** *smt*

theorem *Th4*: $\lfloor \diamond(\exists^E x. G\ x) \rfloor$ **using** *Ax2a Ax4 L Rsymmm Th2 by metis*

theorem *Th5*: $\lfloor \square(\exists^E x. G\ x) \rfloor$ **using** *Th3 Th4 by blast*

lemma *MC*: $\lfloor \varphi \supset \square\varphi \rfloor$

— sledgehammer(*Ax2a Ax2b Th5 God_def Rsymmm*) — proof found

proof — {fix *w* fix *Q*

have 1: $\forall x. (G\ x\ w \longrightarrow (\forall Z. Z\ x \supset \square(\forall z. G\ z \supset Z\ z))\ w)$ **using** *Ax2a Ax2b God-def by smt*

have 2: $(\exists x. G\ x\ w) \longrightarrow ((Q \supset \square(\forall z. G\ z \supset Q))\ w)$ **using** 1 **by force**

have 3: $(Q \supset \square Q)\ w$ **using** 2 *Th5 Rsymmm by blast*

thus *?thesis by auto*

qed

lemma *PosProps*: $\lfloor P\ (\lambda x. \top) \wedge P\ (\lambda x. x = x) \rfloor$ **using** *Ax2a Ax4 by blast*

lemma *NegProps*: $\lfloor \neg P(\lambda x. \perp) \wedge \neg P(\lambda x. x \neq x) \rfloor$ **using** *Ax2a Ax4 by blast*

lemma *UniqueEss1*: $\lfloor \varphi\ Ess.\ x \wedge \psi\ Ess.\ x \supset \square(\forall y. \varphi\ y \leftrightarrow \psi\ y) \rfloor$ **using** *Essence-def by smt*

lemma *UniqueEss2*: $\lfloor \varphi\ Ess.\ x \wedge \psi\ Ess.\ x \supset \square(\varphi \equiv \psi) \rfloor$ **nitpick**[*card i=2*] **oops**
— Countermodel found

lemma *UniqueEss3*: $\lfloor \varphi\ Ess.\ x \supset \square(\forall y. \varphi\ y \supset y \equiv x) \rfloor$ **using** *Essence-def MC*

```

by auto
lemma Monotheism: [ $G x \wedge G y \supset x \equiv y$ ] using Ax2a God-def by smt
lemma Filter: [ $\text{Filter } P$ ] using Ax1 Ax4 MC NegProps PosProps Rsymm Ax2a
God-def Th5 by (smt (verit, ccfv-threshold))
lemma UltraFilter: [ $\text{UFilter } P$ ] using Ax2a Filter by smt
lemma True nitpick[satisfy,card=1,eval= $[P (\lambda x. \perp)]$ ] oops — One model found
of cardinality one

end

```

7.3 GoedelVariantHOML3AndersonQuan.thy

The same as GoedelVariantHOML3, but now for a mixed use of actualist and possibilist quantifiers for entities; cf. Footnote 20 in [2].

```

theory GoedelVariantHOML3AndersonQuant imports HOMLinHOL ModalFilter
begin

```

```
consts PositiveProperty:: $(e \Rightarrow \sigma) \Rightarrow \sigma (P)$ 
```

```
axiomatization where Ax1: [ $P \varphi \wedge P \psi \supset P (\varphi . \psi)$ ]
```

```
axiomatization where Ax2a: [ $P \varphi \vee^e P \sim \varphi$ ]
```

```
definition God (G) where  $G x \equiv \forall \varphi. P \varphi \supset \varphi x$ 
```

```
abbreviation PropertyInclusion ( $\neg\supset_N$ ) where  $\varphi \supset_N \psi \equiv \square(\varphi \neq (\lambda x. \perp) \wedge (\forall y::e. \varphi y \supset \psi y))$ 
```

```
definition Essence (-Ess.-) where  $\varphi \text{ Ess. } x \equiv \forall \psi. \psi x \supset (\varphi \supset_N \psi)$ 
```

```
axiomatization where Ax2b: [ $P \varphi \supset \square P \varphi$ ]
```

```
lemma Ax2b': [ $\neg P \varphi \supset \square(\neg P \varphi)$ ] using Ax2a Ax2b by blast
```

```
theorem Th1: [ $G x \supset G \text{ Ess. } x$ ] using Ax2a Ax2b Essence-def God-def by (smt (verit))
```

```
definition NecExist (E) where  $E x \equiv \forall \varphi. \varphi \text{ Ess. } x \supset \square(\exists^E x. \varphi x)$ 
```

```
axiomatization where Ax3: [ $P E$ ]
```

```
theorem Th2: [ $G x \supset \square(\exists^E y. G y)$ ] using Ax3 Th1 God-def NecExist-def by smt
```

```
theorem Th3: [ $\diamond(\exists^E x. G x) \supset \square(\exists^E y. G y)$ ]
```

— sledgehammer(Th2 Rsymm) — Proof found

proof —

```
have 1: [ $(\exists^E x. G x) \supset \square(\exists^E y. G y)$ ] using Th2 by blast
```

```
have 2: [ $\diamond(\exists^E x. G x) \supset \diamond\square(\exists^E y. G y)$ ] using 1 by blast
```

have 3: $\lfloor \Diamond(\exists^E x. G x) \supset \Box(\exists^E y. G y) \rfloor$ **using** 2 *Rsymm* **by** *blast*
thus ?*thesis* **by** *blast*
qed

axiomatization where *Ax4*: $\lfloor P \varphi \wedge (\varphi \supset_N \psi) \supset P \psi \rfloor$

lemma *True nitpick*[*satisfy*,*card*=1,*eval*= $\lfloor P(\lambda x. \perp) \rfloor$] **oops** — Two models found of cardinality one

abbreviation *PosProps* $\Phi \equiv \forall \varphi. \Phi \varphi \supset P \varphi$
abbreviation *ConjOfPropsFrom* $\varphi \Phi \equiv \Box(\forall^E z. \varphi z \leftrightarrow (\forall \psi. \Phi \psi \supset \psi z))$
axiomatization where *Ax1Gen*: $\lfloor (\text{PosProps } \Phi \wedge \text{ConjOfPropsFrom } \varphi \Phi) \supset P \varphi \rfloor$

lemma *L*: $\lfloor P G \rfloor$ **using** *Ax1Gen* *God-def* **by** (*smt (verit)*)

theorem *Th4*: $\lfloor \Diamond(\exists^E x. G x) \rfloor$
— sledgehammer[timeout=200](*Ax2a L Ax1Gen*) **oops** — sorry — Proof found

axiomatization where *Th4*: $\lfloor \Diamond(\exists^E x. G x) \rfloor$

theorem *Th5*: $\lfloor \Box(\exists^E x. G x) \rfloor$ **using** *Th3 Th4* **by** *blast*

lemma *MC*: $\lfloor \varphi \supset \Box \varphi \rfloor$
— sledgehammer(*Ax2a Ax2b Th5 God_def Rsymm*) — Proof found
proof — {fix *w* fix *Q*
have 1: $\forall x. (G x w \longrightarrow (\forall Z. Z x \supset \Box(\forall z. G z \supset Z z)) w)$ **using** *Ax2a Ax2b God-def* **by** *smt*
have 2: $(\exists x. G x w) \longrightarrow ((Q \supset \Box(\forall z. G z \supset Q)) w)$ **using** 1 **by** *force*
have 3: $(Q \supset \Box Q) w$ **using** 2 *Th5 Rsymm* **by** *blast*}
thus ?*thesis* **by** *auto*
qed

lemma *PosProps*: $\lfloor P(\lambda x. \top) \wedge P(\lambda x. x = x) \rfloor$ **using** *Ax2a Ax4 L Th4* **by** *smt*

lemma *NegProps*: $\lfloor \neg P(\lambda x. \perp) \wedge \neg P(\lambda x. x \neq x) \rfloor$ **using** *Ax2a Ax4 L Th4* **by** *smt*

lemma *UniqueEss1*: $\lfloor \varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \Box(\forall y. \varphi y \leftrightarrow \psi y) \rfloor$ **oops** — Unclear, open question

lemma *UniqueEss2*: $\lfloor \varphi \text{ Ess. } x \wedge \psi \text{ Ess. } x \supset \Box(\varphi \equiv \psi) \rfloor$ **oops** — Unclear, open question

lemma *UniqueEss3*: $\lfloor \varphi \text{ Ess. } x \supset \Box(\forall y. \varphi y \supset y \equiv x) \rfloor$ **using** *Essence-def MC by auto*

lemma *Monotheism*: $\lfloor G x \wedge G y \supset x \equiv y \rfloor$ **using** *Ax2a God-def* **by** *smt*

lemma *Filter*: $\lfloor \text{Filter } P \rfloor$ **using** *Ax1 Ax4 MC NegProps PosProps Rsymm Ax2a God-def Th5* **by** (*smt (verit, ccfv-threshold)*)

lemma *UltraFilter*: $\lfloor \text{UFilter } P \rfloor$ **using** *Ax2a Filter* **by** *blast*

lemma *True nitpick*[*satisfy*,*card*=1,*eval*= $\lfloor P(\lambda x. \top) \rfloor$] **oops** — One model found of cardinality one

end

7.4 HOMLinHOLonlyS4.thy (slight variation of Figure 3 of [2])

Shallow embedding of higher-order modal logic (HOML) in the classical higher-order logic (HOL) of Isabelle/HOL utilizing the LogiKEy methodology. Here logic S4 is introduced.

```
theory HOMLinHOLonlyS4 imports Main
begin
```

— Global parameters setting for the model finder nitpick and the parser; unimport for the reader

```
nitpick_params[user-axioms,expect=genuine,show-all,format=2,max-genuine=3]
declare[[syntax-ambiguity-warning=false]]
```

— Type i is associated with possible worlds and type e with entities:

```
typedecl i — Possible worlds
```

```
typedecl e — Individuals/entities
```

```
type-synonym  $\sigma = i \Rightarrow \text{bool}$  — World-lifted propositions
```

```
type-synonym  $\tau = e \Rightarrow \sigma$  — modal properties
```

```
consts R:: $i \Rightarrow i \Rightarrow \text{bool}$  (-r-) — Accessibility relation between worlds
```

```
axiomatization where
```

```
Rrefl:  $\forall x. xrx$  and
```

```
Rtrans:  $\forall x y z. xry \wedge yrz \longrightarrow xrz$ 
```

— Logical connectives (operating on truth-sets)

```
abbreviation Mbot:: $\sigma$  ( $\perp$ ) where  $\perp \equiv \lambda w. \text{False}$ 
```

```
abbreviation Mtop:: $\sigma$  ( $\top$ ) where  $\top \equiv \lambda w. \text{True}$ 
```

```
abbreviation Mneg:: $\sigma \Rightarrow \sigma$  ( $\neg$ - [52]53) where  $\neg\varphi \equiv \lambda w. \neg(\varphi w)$ 
```

```
abbreviation Mand:: $\sigma \Rightarrow \sigma \Rightarrow \sigma$  (infixl  $\wedge$  50) where  $\varphi \wedge \psi \equiv \lambda w. \varphi w \wedge \psi w$ 
```

```
abbreviation Mor:: $\sigma \Rightarrow \sigma \Rightarrow \sigma$  (infixl  $\vee$  49) where  $\varphi \vee \psi \equiv \lambda w. \varphi w \vee \psi w$ 
```

```
abbreviation Mimp:: $\sigma \Rightarrow \sigma \Rightarrow \sigma$  (infixr  $\supset$  48) where  $\varphi \supset \psi \equiv \lambda w. \varphi w \longrightarrow \psi w$ 
```

```
abbreviation Mequiv:: $\sigma \Rightarrow \sigma \Rightarrow \sigma$  (infixl  $\leftrightarrow$  47) where  $\varphi \leftrightarrow \psi \equiv \lambda w. \varphi w \longleftrightarrow \psi w$ 
```

```
abbreviation Mbox:: $\sigma \Rightarrow \sigma$  ( $\Box$ - [54]55) where  $\Box \varphi \equiv \lambda w. \forall v. w \mathbf{r} v \longrightarrow \varphi v$ 
```

```
abbreviation Mdia:: $\sigma \Rightarrow \sigma$  ( $\Diamond$ - [54]55) where  $\Diamond \varphi \equiv \lambda w. \exists v. w \mathbf{r} v \wedge \varphi v$ 
```

```
abbreviation Mprimeq::' $a \Rightarrow a \Rightarrow \sigma$  ( $=\!=$ ) where  $x=y \equiv \lambda w. x=y$ 
```

```
abbreviation Mprimneg::' $a \Rightarrow a \Rightarrow \sigma$  ( $\neq$ ) where  $x \neq y \equiv \lambda w. x \neq y$ 
```

```
abbreviation Mnegpred:: $\tau \Rightarrow \tau$  ( $\sim$ -) where  $\sim \Phi \equiv \lambda x. \lambda w. \neg \Phi x w$ 
```

```
abbreviation Mconpred:: $\tau \Rightarrow \tau \Rightarrow \tau$  (infixl . 50) where  $\Phi. \Psi \equiv \lambda x. \lambda w. \Phi x w \wedge \Psi x w$ 
```

```
abbreviation Mexclor:: $\sigma \Rightarrow \sigma \Rightarrow \sigma$  (infixl  $\vee^e$  49) where  $\varphi \vee^e \psi \equiv (\varphi \vee \psi) \wedge \neg(\varphi \wedge \psi)$ 
```

— Possibilist quantifiers (polymorphic)

```
abbreviation Mallposs::(' $a \Rightarrow \sigma \Rightarrow \sigma$ ) ( $\forall$ ) where  $\forall \Phi \equiv \lambda w. \forall x. \Phi x w$ 
```

```
abbreviation Mallpossb (binder  $\forall$  [8]9) where  $\forall x. \varphi(x) \equiv \forall \varphi$ 
```

```
abbreviation Mexiposs::(' $a \Rightarrow \sigma \Rightarrow \sigma$ ) ( $\exists$ ) where  $\exists \Phi \equiv \lambda w. \exists x. \Phi x w$ 
```

```
abbreviation Mexipossb (binder  $\exists$  [8]9) where  $\exists x. \varphi(x) \equiv \exists \varphi$ 
```

- Actualist quantifiers (for individuals/entities)


```
consts existsAt:: $e \Rightarrow \sigma$  (-@-)
abbreviation Mallact:: $(e \Rightarrow \sigma) \Rightarrow \sigma$  ( $\forall^E$ ) where  $\forall^E \Phi \equiv \lambda w. \forall x. x @ w \rightarrow \Phi x w$ 
abbreviation Mallactb (binder  $\forall^E$  [8]9) where  $\forall^E x. \varphi(x) \equiv \forall^E \varphi$ 
abbreviation Mexiact:: $(e \Rightarrow \sigma) \Rightarrow \sigma$  ( $\exists^E$ ) where  $\exists^E \Phi \equiv \lambda w. \exists x. x @ w \wedge \Phi x w$ 
abbreviation Mexiactb (binder  $\exists^E$  [8]9) where  $\exists^E x. \varphi(x) \equiv \exists^E \varphi$ 
```
- Leibniz equality (polymorphic)


```
abbreviation Mleibeq::' $a \Rightarrow a \Rightarrow \sigma$ ' (-≡-) where  $x \equiv y \equiv \forall P. P x \supset P y$ 
```
- Meta-logical predicate for global validity


```
abbreviation Mvalid:: $\sigma \Rightarrow \text{bool}$  ([ - ]) where  $[\psi] \equiv \forall w. \psi w$ 
```

end

7.5 TestsHOMLinS4.thy

Tests and verifications of properties for the embedding of HOML (S4) in HOL.

```
theory TestsHOMLinS4 imports HOMLinHOLonlyS4
begin
```

- Test for S5 modal logic


```
lemma axM:  $[\Box\varphi \supset \varphi]$  using Rrefl by blast
lemma axD:  $[\Box\varphi \supset \Diamond\varphi]$  using Rrefl by blast
lemma axB:  $[\varphi \supset \Box\Diamond\varphi]$ 
nitpick[expect=genuine] oops — Countermodel found
lemma ax4:  $[\Box\varphi \supset \Box\Box\varphi]$  using Rtrans by blast
lemma ax5:  $[\Diamond\varphi \supset \Box\Diamond\varphi]$ 
nitpick[expect=genuine] oops — Countermodel found
```

- Test for Barcan and converse Barcan formula:


```
lemma BarcanAct:  $[(\forall^E x. \Box(\varphi x)) \supset \Box(\forall^E x. (\varphi x))]$ 
nitpick[expect=genuine] oops — Countermodel found
lemma ConvBarcanAct:  $[\Box(\forall^E x. (\varphi x)) \supset (\forall^E x. \Box(\varphi x))]$ 
nitpick[expect=genuine] oops — Countermodel found
lemma BarcanPoss:  $[(\forall x. \Box(\varphi x)) \supset \Box(\forall x. \varphi x)]$  by blast
lemma ConvBarcanPoss:  $[\Box(\forall x. (\varphi x)) \supset (\forall x. \Box(\varphi x))]$  by blast
```

- A simple Hilbert system for classical propositional logic is derived


```
lemma Hilbert-A1:  $[A \supset (B \supset A)]$  by blast
lemma Hilbert-A2:  $[(A \supset (B \supset C)) \supset ((A \supset B) \supset (A \supset C))]$  by blast
lemma Hilbert-MP: assumes  $[A]$  and  $[A \supset B]$  shows  $[B]$  using assms by blast
```

- We have a polymorphic possibilist quantifier for which existential import holds


```
lemma Quant-1: assumes  $[A]$  shows  $[\forall x::'a. A]$  using assms by auto
```

- Existential import holds for possibilist quantifiers

```

lemma ExImPossibilist1:  $\exists x::e. x = x$  by blast
lemma ExImPossibilist2:  $\exists x::e. x \equiv x$  by blast
lemma ExImPossibilist3:  $\exists x::e. x = t$  by blast
lemma ExImPossibilist4:  $\exists x::'a. x \equiv t::'a$  by blast
lemma ExImPossibilist:  $\exists x::'a. \top$  by blast

```

— We have an actualist quantifier for individuals for which existential import does not hold

```
lemma Quant-2: assumes  $[A]$  shows  $[\forall^E x::e. A]$  using assms by auto
```

— Existential import does not hold for our actualist quantifiers (for individuals)

```

lemma ExImActualist1:  $\exists^E x::e. x = x$ 
  nitpick[card=1,expect=genuine] oops — Countermodel found
lemma ExImActualist2:  $\exists^E x::e. x \equiv x$ 
  nitpick[card=1,expect=genuine] oops — Countermodel found
lemma ExImActualist3:  $\exists^E x::e. x = t$ 
  nitpick[card=1,expect=genuine] oops — Countermodel found
lemma ExImActualist:  $\exists^E x::e. \top$ 
  nitpick[card=1,expect=genuine] oops — Countermodel found

```

— Properties of the embedded primitive equality, which coincides with Leibniz equality

```

lemma EqRefl:  $[x = x]$  by blast
lemma EqSym:  $[(x = y) \leftrightarrow (y = x)]$  by blast
lemma EqTrans:  $[(x = y) \wedge (y = z) \supset (x = z)]$  by blast
lemma EQCong:  $[(x = y) \supset ((\varphi x) = (\varphi y))]$  by blast
lemma EQFuncExt:  $[(\varphi = \psi) \supset (\forall x. ((\varphi x) = (\psi x)))]$  by blast
lemma EQBoolExt1:  $[(\varphi = \psi) \supset (\varphi \leftrightarrow \psi)]$  by blast
lemma EQBoolExt2:  $[(\varphi \leftrightarrow \psi) \supset (\varphi = \psi)]$ 
  nitpick[card=2] oops — Countermodel found
lemma EQBoolExt3:  $[(\varphi \leftrightarrow \psi)] \longrightarrow [(\varphi = \psi)]$  by blast
lemma EqPrimLeib:  $[(x = y) \leftrightarrow (x \equiv y)]$  by auto

```

— Comprehension is natively supported in HOL (due to lambda-abstraction)

```

lemma Comprehension1:  $[\exists \varphi. \forall x. (\varphi x) \leftrightarrow A]$  by force
lemma Comprehension2:  $[\exists \varphi. \forall x. (\varphi x) \leftrightarrow (A x)]$  by force
lemma Comprehension3:  $[\exists \varphi. \forall x y. (\varphi x y) \leftrightarrow (A x y)]$  by force

```

— Modal collapse does not hold

```
lemma ModalCollapse:  $[\forall \varphi. \varphi \supset \square \varphi]$ 
  nitpick[card=2,expect=genuine] oops — Countermodel found
```

— Empty property and self-difference

```

lemma TruePropertyAndSelfIdentity:  $[(\lambda x::e. \top) = (\lambda x. x = x)]$  by blast
lemma EmptyPropertyAndSelfDifference:  $[(\lambda x::e. \perp) = (\lambda x. x \neq x)]$  by blast
lemma EmptyProperty2:  $[\exists x. \varphi x] \longrightarrow [\varphi \neq (\lambda x::e. \perp)]$  by blast
lemma EmptyProperty3:  $[\exists^E x. \varphi x] \longrightarrow [\varphi \neq (\lambda x::e. \perp)]$  by blast
lemma EmptyProperty4:  $[\varphi \neq (\lambda x::e. \perp)] \longrightarrow [\exists x. \varphi x]$ 
  nitpick[expect=genuine] oops — Countermodel found

```

```

lemma EmptyProperty5: [ $\varphi \neq (\lambda x::e. \perp)$ ] —> [ $\exists^E x. \varphi x$ ]
nitpick[expect=genuine] oops — Countermodel found
end

```

7.6 GoedelVariantHOML1inS4.thy

The same as GoedelVariantHOML1, but now in logic S4.

```

theory GoedelVariantHOML1inS4 imports HOMLInHOLonlyS4
begin

```

```

consts PositiveProperty::( $e \Rightarrow \sigma \Rightarrow \sigma$ ) ( $P$ )

```

```

axiomatization where Ax1: [ $P \varphi \wedge P \psi \supset P(\varphi \cdot \psi)$ ]

```

```

axiomatization where Ax2a: [ $P \varphi \vee^e P \sim \varphi$ ]

```

```

definition God (G) where  $G x \equiv \forall \varphi. P \varphi \supset \varphi x$ 

```

```

abbreviation PropertyInclusion (- $\supset_N$ -) where  $\varphi \supset_N \psi \equiv \square(\forall^E y. \varphi y \supset \psi y)$ 

```

```

definition Essence (- $\text{Ess.}$ -) where  $\varphi \text{ Ess. } x \equiv \forall \psi. \psi x \supset (\varphi \supset_N \psi)$ 

```

```

axiomatization where Ax2b: [ $P \varphi \supset \square P \varphi$ ]

```

```

lemma Ax2b': [ $\neg P \varphi \supset \square(\neg P \varphi)$ ] using Ax2a Ax2b by blast

```

```

theorem Th1: [ $G x \supset G \text{ Ess. } x$ ] using Ax2a Ax2b Essence-def God-def by (smt (verit))

```

```

definition NecExist (E) where  $E x \equiv \forall \varphi. \varphi \text{ Ess. } x \supset \square(\exists^E x. \varphi x)$ 

```

```

axiomatization where Ax3: [ $P E$ ]

```

```

theorem Th2: [ $G x \supset \square(\exists^E y. G y)$ ] using Ax3 Th1 God-def NecExist-def by
smt

```

```

axiomatization where Ax4: [ $P \varphi \wedge (\varphi \supset_N \psi) \supset P \psi$ ]

```

```

theorem Th3: [ $\diamond(\exists^E x. G x) \supset \square(\exists^E y. G y)$ ] using Ax3 Essence-def God-def
NecExist-def Rrefl by fastforce

```

```

end

```

7.7 GoedelVariantHOML2inS4.thy

The same as GoedelVariantHOML2, but now in logic S4, where the proof of theorem Th3 fails.

```

theory GoedelVariantHOML2inS4 imports HOMLinHOLonlyS4
begin

consts PositiveProperty::(e⇒σ)⇒σ (P)

axiomatization where Ax1: [P φ ∧ P ψ ⊃ P (φ ∙ ψ)] 

abbreviation PosProps Φ ≡ ∀φ. Φ φ ⊃ P φ
abbreviation ConjOfPropsFrom φ Φ ≡ □(∀Ez. φ z ↔ (∀ψ. Φ ψ ⊃ ψ z))
axiomatization where Ax1Gen: [(PosProps Φ ∧ ConjOfPropsFrom φ Φ) ⊃ P φ]

axiomatization where Ax2a: [P φ ∨e P ∼φ] 

definition God (G) where G x ≡ ∀φ. P φ ⊃ φ x

abbreviation PropertyInclusion (-▷N-) where φ ▷N ψ ≡ □(∀Ey. φ y ⊃ ψ y)

definition Essence (-Ess.-) where φ Ess. x ≡ φ x ∧ (∀ψ. ψ x ⊃ (φ ▷N ψ))

axiomatization where Ax2b: [P φ ⊃ □ P φ]

lemma Ax2b': [¬P φ ⊃ □(¬P φ)] using Ax2a Ax2b by blast

theorem Th1: [G x ⊃ G Ess. x] using Ax2a Ax2b Essence-def God-def by (smt (verit))

definition NecExist (E) where E x ≡ ∀φ. φ Ess. x ⊃ □(∃Ex. φ x)

axiomatization where Ax3: [P E]

axiomatization where Ax4: [P φ ∧ (φ ▷N ψ) ⊃ P ψ]

theorem Th2: [G x ⊃ □(∃Ey. G y)] using Ax3 Th1 God-def NecExist-def by smt

theorem Th3: [◊(∃Ex. G x) ⊃ □(∃Ey. G y)] — nitpick sledgehammer oops
— Open problem

end

```

7.8 GoedelVariantHOML2possInS4.thy

The same as GoedelVariantHOML2poss, but now in logic S4, where the proof of theorem Th3 fails.

```

theory GoedelVariantHOML2possInS4 imports HOMLinHOLonlyS4
begin

```

```

consts PositiveProperty::(e⇒σ)⇒σ (P)

```

axiomatization where *Ax1*: $[P \varphi \wedge P \psi \supset P (\varphi \cdot \psi)]$

abbreviation *PosProps* $\Phi \equiv \forall \varphi. \Phi \varphi \supset P \varphi$

abbreviation *ConjOfPropsFrom* $\varphi \Phi \equiv \square(\forall z. \varphi z \leftrightarrow (\forall \psi. \Phi \psi \supset \psi z))$

axiomatization where *Ax1Gen*: $[(\text{PosProps } \Phi \wedge \text{ConjOfPropsFrom } \varphi \Phi) \supset P \varphi]$

axiomatization where *Ax2a*: $[P \varphi \vee^e P \sim \varphi]$

definition *God (G)* **where** $G x \equiv \forall \varphi. P \varphi \supset \varphi x$

abbreviation *PropertyInclusion (- \supset_N -)* **where** $\varphi \supset_N \psi \equiv \square(\forall y::e. \varphi y \supset \psi y)$

definition *Essence (- Ess. -)* **where** $\varphi \text{ Ess. } x \equiv \varphi x \wedge (\forall \psi. \psi x \supset (\varphi \supset_N \psi))$

axiomatization where *Ax2b*: $[P \varphi \supset \square P \varphi]$

lemma *Ax2b'*: $[\neg P \varphi \supset \square(\neg P \varphi)]$ **using** *Ax2a Ax2b* **by** *blast*

theorem *Th1*: $[G x \supset G \text{ Ess. } x]$ **using** *Ax2a Ax2b Essence-def God-def* **by** (*smt (verit)*)

definition *NecExist (E)* **where** $E x \equiv \forall \varphi. \varphi \text{ Ess. } x \supset \square(\exists x. \varphi x)$

axiomatization where *Ax3*: $[P E]$

axiomatization where *Ax4*: $[P \varphi \wedge (\varphi \supset_N \psi) \supset P \psi]$

theorem *Th2*: $[G x \supset \square(\exists y. G y)]$ **using** *Ax3 Th1 God-def NecExist-def* **by** *smt*

theorem *Th3*: $[\diamond(\exists x. G x) \supset \square(\exists y. G y)]$ — nitpick sledgehammer **oops** — Open problem

end

7.9 GoedelVariantHOML3inS4.thy

The same as GoedelVariantHOML3, but now in logic S4, where the proof of theorem Th3 fails.

theory *GoedelVariantHOML3inS4* **imports** *HOMLinHOLonlyS4*
begin

consts *PositiveProperty*:: $(e \Rightarrow \sigma) \Rightarrow \sigma (P)$

axiomatization where *Ax1*: $[P \varphi \wedge P \psi \supset P (\varphi \cdot \psi)]$

abbreviation *PosProps* $\Phi \equiv \forall \varphi. \Phi \varphi \supset P \varphi$
abbreviation *ConjOfPropsFrom* $\varphi \Phi \equiv \square(\forall^E z. \varphi z \leftrightarrow (\forall \psi. \Phi \psi \supset \psi z))$
axiomatization where *Ax1Gen*: $\lfloor (\text{PosProps } \Phi \wedge \text{ConjOfPropsFrom } \varphi \Phi) \supset P \varphi \rfloor$

axiomatization where *Ax2a*: $\lfloor P \varphi \vee^e P \sim \varphi \rfloor$

definition *God* (*G*) **where** $G x \equiv \forall \varphi. P \varphi \supset \varphi x$

abbreviation *PropertyInclusion* ($\neg \supset_N$) **where** $\varphi \supset_N \psi \equiv \square(\varphi \neq (\lambda x. \perp) \wedge (\forall^E y. \varphi y \supset \psi y))$

definition *Essence* ($\neg \text{Ess}.$) **where** $\varphi \text{ Ess. } x \equiv \forall \psi. \psi x \supset (\varphi \supset_N \psi)$

axiomatization where *Ax2b*: $\lfloor P \varphi \supset \square P \varphi \rfloor$

lemma *Ax2b'*: $\lfloor \neg P \varphi \supset \square(\neg P \varphi) \rfloor$ **using** *Ax2a Ax2b* **by** *blast*

theorem *Th1*: $\lfloor G x \supset G \text{ Ess. } x \rfloor$ **using** *Ax2a Ax2b Essence-def God-def* **by** (*smt (verit)*)

definition *NecExist* (*E*) **where** $E x \equiv \forall \varphi. \varphi \text{ Ess. } x \supset \square(\exists^E x. \varphi x)$

axiomatization where *Ax3*: $\lfloor P E \rfloor$

theorem *Th2*: $\lfloor G x \supset \square(\exists^E y. G y) \rfloor$ **using** *Ax3 Th1 God-def NecExist-def* **by** *smt*

axiomatization where *Ax4*: $\lfloor P \varphi \wedge (\varphi \supset_N \psi) \supset P \psi \rfloor$

theorem *Th3*: $\lfloor \diamond(\exists^E x. G x) \supset \square(\exists^E y. G y) \rfloor$ — nitpick sledgehammer **oops**
— Open problem

end

7.10 GoedelVariantHOML3possInS4.thy

The same as GoedelVariantHOML3poss, but now in logic S4, where the proof of theorem Th3 fails.

theory *GoedelVariantHOML3possInS4* **imports** *HOMLinHOLonlyS4*
begin

consts *PositiveProperty*:: $(e \Rightarrow \sigma) \Rightarrow \sigma$ (*P*)

axiomatization where *Ax1*: $\lfloor P \varphi \wedge P \psi \supset P (\varphi \cdot \psi) \rfloor$

abbreviation *PosProps* $\Phi \equiv \forall \varphi. \Phi \varphi \supset P \varphi$

abbreviation *ConjOfPropsFrom* $\varphi \Phi \equiv \square(\forall z. \varphi z \leftrightarrow (\forall \psi. \Phi \psi \supset \psi z))$

axiomatization where *Ax1Gen*: $\lfloor (PosProps \Phi \wedge ConjOfPropsFrom \varphi \Phi) \supset P \varphi \rfloor$

axiomatization where *Ax2a*: $\lfloor P \varphi \vee^e P \sim \varphi \rfloor$

definition *God* (*G*) **where** $G x \equiv \forall \varphi. P \varphi \supset \varphi x$

abbreviation *PropertyInclusion* ($\neg\supset_N$) **where** $\varphi \supset_N \psi \equiv \square(\varphi \neq (\lambda x. \perp) \wedge (\forall y. \varphi y \supset \psi y))$

definition *Essence* ($\neg\text{Ess.}$) **where** $\varphi \text{ Ess. } x \equiv \forall \psi. \psi x \supset (\varphi \supset_N \psi)$

axiomatization where *Ax2b*: $\lfloor P \varphi \supset \square P \varphi \rfloor$

lemma *Ax2b'*: $\lfloor \neg P \varphi \supset \square(\neg P \varphi) \rfloor$ **using** *Ax2a Ax2b* **by** *blast*

theorem *Th1*: $\lfloor G x \supset G \text{ Ess. } x \rfloor$ **using** *Ax2a Ax2b Essence-def God-def* **by** (*smt (verit)*)

definition *NecExist* (*E*) **where** $E x \equiv \forall \varphi. (\varphi \text{ Ess. } x) \supset \square(\exists x. \varphi x)$

axiomatization where *Ax3*: $\lfloor P E \rfloor$

axiomatization where *Ax4*: $\lfloor P \varphi \wedge (\varphi \supset_N \psi) \supset P \psi \rfloor$

theorem *Th2*: $\lfloor G x \supset \square(\exists y. G y) \rfloor$ **using** *Ax3 Th1 God-def NecExist-def* **by** *smt*

theorem *Th3*: $\lfloor \diamond(\exists x. G x) \supset \square(\exists y. G y) \rfloor$ — nitpick sledgehammer **oops** — Open problem

end

7.11 ScottVariantHOMLinS4.thy

theory *ScottVariantHOMLinS4* **imports** *HOMLinHOLonlyS4*
begin

consts *PositiveProperty*:: $(e \Rightarrow \sigma) \Rightarrow \sigma$ (*P*)

axiomatization where *A1*: $\lfloor \neg P \varphi \leftrightarrow P \sim \varphi \rfloor$

axiomatization where *A2*: $\lfloor P \varphi \wedge \square(\forall^E y. \varphi y \supset \psi y) \supset P \psi \rfloor$

theorem *T1*: $\lfloor P \varphi \supset \diamond(\exists^E x. \varphi x) \rfloor$ **using** *A1 A2* **by** *blast*

definition *God* (*G*) **where** $G x \equiv \forall \varphi. P \varphi \supset \varphi x$

axiomatization where *A3*: $\lfloor P G \rfloor$

theorem *Coro*: $\lfloor \Diamond(\exists^E x. G x) \rfloor$ **using** *A3 T1* **by** *blast*

axiomatization where *A4*: $\lfloor P \varphi \supset \Box P \varphi \rfloor$

definition *Ess (-Ess.-)* **where** φ *Ess.* $x \equiv \varphi x \wedge (\forall \psi. \psi x \supset \Box(\forall^E y. \varphi y \supset \psi y))$

theorem *T2*: $\lfloor G x \supset G \text{Ess. } x \rfloor$ **using** *A1 A4 Ess-def God-def* **by** *fastforce*

definition *NecExist (NE)* **where** $NE x \equiv \forall \varphi. \varphi \text{Ess. } x \supset \Box(\exists^E x. \varphi x)$

axiomatization where *A5*: $\lfloor P NE \rfloor$

lemma *True nitpick*[*satisfy, card=1, eval=[P (λx.T)]*] **oops** — One model found of cardinality one

theorem *T3*: $\lfloor \Box(\exists^E x. G x) \rfloor$ **nitpick**[*card e=1, card i=2*] **oops** — Countermodel found

lemma *MC*: $\lfloor \varphi \supset \Box \varphi \rfloor$ **nitpick**[*card e=1, card i=2*] **oops** — Countermodel found

end

7.12 ScottVariantHOMLpossInS4.thy

theory *ScottVariantHOMLpossInS4 imports HOMLinHOLonlyS4*
begin

consts *PositiveProperty::(e⇒σ)⇒σ (P)*

axiomatization where *A1*: $\lfloor \neg P \varphi \leftrightarrow P \sim \varphi \rfloor$

axiomatization where *A2*: $\lfloor P \varphi \wedge \Box(\forall y. \varphi y \supset \psi y) \supset P \psi \rfloor$

theorem *T1*: $\lfloor P \varphi \supset \Diamond(\exists x. \varphi x) \rfloor$ **using** *A1 A2* **by** *blast*

definition *God (G)* **where** $G x \equiv \forall \varphi. P \varphi \supset \varphi x$

axiomatization where *A3*: $\lfloor P G \rfloor$

theorem *Coro*: $\lfloor \Diamond(\exists x. G x) \rfloor$ **using** *A3 T1* **by** *blast*

axiomatization where *A4*: $\lfloor P \varphi \supset \Box P \varphi \rfloor$

definition *Ess (-Ess.-)* **where** φ *Ess.* $x \equiv \varphi x \wedge (\forall \psi. \psi x \supset \Box(\forall y::e. \varphi y \supset \psi y))$

theorem *T2*: $\lfloor G x \supset G \text{Ess. } x \rfloor$ **using** *A1 A4 Ess-def God-def* **by** *fastforce*

```

definition NecExist (NE) where NE x ≡ ∀ φ. φ Ess. x ⊃ □(∃ x. φ x)

axiomatization where A5: [P NE]

lemma True nitpick[satisfy,card=1,eval=[P (λx.⊥)]] oops — One model found
of cardinality one

theorem T3: [□(∃ x. G x)] nitpick[card e=1, card i=2] oops — Countermodel
found

lemma MC: [φ ⊃ □φ] nitpick[card e=1, card i=2] oops — Countermodel found

end

```

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