

The Localization of a Commutative Ring

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Abstract

We formalize the localization [1, II, §4] of a commutative ring R with respect to a multiplicative subset (i.e. a submonoid of R seen as a multiplicative monoid).

This localization is itself a commutative ring and we build the natural homomorphism of rings from R to its localization.

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theory *Localization*

imports *Main HOL-Algebra.Group HOL-Algebra.Ring HOL-Algebra.AbelCoset*
begin

Contents:

- We define the localization of a commutative ring R with respect to a multiplicative subset, i.e. with respect to a submonoid of R (seen as a multiplicative monoid), cf. [*rec-rng-of-frac*].
- We prove that this localization is a commutative ring (cf. [*crng-rng-of-frac*]) equipped with a homomorphism of rings from R (cf. [*rng-to-rng-of-frac-is-ring-hom*]).

1 The Localization of a Commutative Ring

1.1 Localization

locale *submonoid = monoid M* **for** M (**structure**) +

fixes S

assumes *subset* : $S \subseteq \text{carrier } M$

and *m-closed* [*intro*, *simp*] : $\llbracket x \in S; y \in S \rrbracket \implies x \otimes y \in S$
and *one-closed* [*simp*] : $\mathbf{1} \in S$

lemma (**in** *submonoid*) *is-submonoid*: *submonoid* $M S$
by (*rule submonoid-axioms*)

locale *mult-submonoid-of-rng* = *ring* $R + \text{submonoid } R S$ **for** R **and** S

locale *mult-submonoid-of-crng* = *cring* $R + \text{mult-submonoid-of-rng } R S$ **for** R
and S

locale *eq-obj-rng-of-frac* = *cring* $R + \text{mult-submonoid-of-crng } R S$ **for** R (**structure**)
and $S +$
fixes *rel*
defines $rel \equiv (\text{carrier} = \text{carrier } R \times S, eq = \lambda(r,s) (r',s'). \exists t \in S. t \otimes ((s' \otimes r) \ominus (s \otimes r')) = \mathbf{0})$

lemma (**in** *abelian-group*) *minus-to-eq* :
assumes *abelian-group* G **and** $x \in \text{carrier } G$ **and** $y \in \text{carrier } G$ **and** $x \ominus y = \mathbf{0}$
shows $x = y$
by (*metis add.inv-solve-right assms(2) assms(3) assms(4) l-zero minus-eq zero-closed*)

lemma (**in** *eq-obj-rng-of-frac*) *equiv-obj-rng-of-frac*:
shows *equivalence* *rel*

proof

show $\bigwedge x. x \in \text{carrier } rel \implies x \text{.}_{rel} x$

proof–

fix x

assume $x \in \text{carrier } rel$

then have $f1: \mathbf{1} \otimes ((snd\ x \otimes fst\ x) \ominus (snd\ x \otimes fst\ x)) = \mathbf{0}$

using *rel-def subset l-one minus-eq add.r-inv rev-subsetD*

by *auto*

moreover have $x = (fst\ x, snd\ x)$

by *simp*

thus $x \text{.}_{rel} x$

using *rel-def one-closed f1*

by *auto*

qed

show $\bigwedge x\ y. x \text{.}_{rel} y \implies x \in \text{carrier } rel \implies y \in \text{carrier } rel \implies y \text{.}_{rel} x$

proof–

fix $x\ y$

assume $a1: x \text{.}_{rel} y$ **and** $a2: x \in \text{carrier } rel$ **and** $a3: y \in \text{carrier } rel$

then obtain t **where** $f1: t \in S$ **and** $f2: t \otimes ((snd\ y \otimes fst\ x) \ominus (snd\ x \otimes fst\ y))$

= $\mathbf{0}$

using *rel-def*

by *fastforce*

then have $(snd\ x \otimes fst\ y) \ominus (snd\ y \otimes fst\ x) = \ominus ((snd\ y \otimes fst\ x) \ominus (snd\ x \otimes fst\ y))$

fst y)

using *abelian-group.minus-add abelian-group.minus-minus*

```

    by (smt a2 a3 a-minus-def abelian-group.a-inv-closed add.inv-mult-group
is-abelian-group
    mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1)
prod.collapse
    rel-def rev-subsetD subset)
  then have  $t \otimes ((snd\ x \otimes fst\ y) \ominus (snd\ y \otimes fst\ x)) = \mathbf{0}$ 
  using minus-zero r-minus f2
  by (smt a2 a3 f1 mem-Sigma-iff minus-closed partial-object.select-convs(1)
prod.collapse
    rel-def semiring-simprules(3) rev-subsetD subset)
  thus  $y \text{ .}=_rel\ x$ 
  using f1 rel-def
  by auto
qed
show  $\bigwedge x\ y\ z.$ 
 $x \text{ .}=_rel\ y \implies y \text{ .}=_rel\ z \implies x \in carrier\ rel \implies y \in carrier\ rel \implies z \in carrier\ rel$ 
 $rel \implies x \text{ .}=_rel\ z$ 
proof-
  fix  $x\ y\ z$ 
  assume  $a1:x \text{ .}=_rel\ y$  and  $a2:y \text{ .}=_rel\ z$  and  $a3:x \in carrier\ rel$  and  $a4:y \in carrier\ rel$ 
  and  $a5:z \in carrier\ rel$ 
  then obtain  $t$  where  $f1:t \in S$  and  $f2:t \otimes ((snd\ y \otimes fst\ x) \ominus (snd\ x \otimes fst\ y)) = \mathbf{0}$ 
  using rel-def
  by fastforce
  then obtain  $t'$  where  $f3:t' \in S$  and  $f4:t' \otimes ((snd\ z \otimes fst\ y) \ominus (snd\ y \otimes fst\ z)) = \mathbf{0}$ 
  using rel-def a2
  by fastforce
  then have  $t \otimes (snd\ y \otimes fst\ x) \ominus t \otimes (snd\ x \otimes fst\ y) = \mathbf{0}$ 
  using f1 subset r-distr f2
  by (smt a3 a4 a-minus-def abelian-group.a-inv-closed is-abelian-group mem-Sigma-iff
    monoid.m-closed monoid-axioms partial-object.select-convs(1) prod.collapse
r-minus rel-def
    subset-iff)
  then have  $t' \otimes (t \otimes (snd\ y \otimes fst\ x)) \ominus t' \otimes (t \otimes (snd\ x \otimes fst\ y)) = \mathbf{0}$ 
  using f3 subset r-distr
  by (smt a3 a4 a-minus-def f1 is-abelian-group mem-Sigma-iff minus-to-eq
    partial-object.select-convs(1) prod.collapse r-neg rel-def semiring-simprules(3)
subset-iff)
  then have  $f5:snd\ z \otimes (t' \otimes (t \otimes (snd\ y \otimes fst\ x))) \ominus snd\ z \otimes (t' \otimes (t \otimes (snd\ x \otimes fst\ y))) = \mathbf{0}$ 
  using a5 rel-def r-distr
  by (smt a3 a4 a-minus-def f1 f3 is-abelian-group mem-Sigma-iff minus-to-eq
    monoid.m-closed
    monoid-axioms partial-object.select-convs(1) prod.collapse r-neg subset
subset-iff)

```

have $t' \otimes (\text{snd } z \otimes \text{fst } y) \ominus t' \otimes (\text{snd } y \otimes \text{fst } z) = \mathbf{0}$
using $f3\ f4\ \text{subset}\ r\text{-distr}$
by (*smt a4 a5 a-minus-def abelian-group.a-inv-closed is-abelian-group mem-Sigma-iff*

monoid.m-closed monoid-axioms partial-object.select-convs(1) prod.collapse
r-minus rel-def
rev-subsetD)

then have $t \otimes (t' \otimes (\text{snd } z \otimes \text{fst } y)) \ominus t \otimes (t' \otimes (\text{snd } y \otimes \text{fst } z)) = \mathbf{0}$
using $f1\ \text{subset}\ r\text{-distr}$
by (*smt a4 a5 a-minus-def f3 is-abelian-group mem-Sigma-iff minus-to-eq*
monoid.m-closed

monoid-axioms partial-object.select-convs(1) prod.collapse r-neg rel-def
subset-iff)

then have $f6:\text{snd } x \otimes (t \otimes (t' \otimes (\text{snd } z \otimes \text{fst } y))) \ominus \text{snd } x \otimes (t \otimes (t' \otimes (\text{snd } y \otimes \text{fst } z))) = \mathbf{0}$
using $a3\ \text{rel-def}\ r\text{-distr}$
by (*smt a4 a5 a-minus-def f1 f3 is-abelian-group mem-Sigma-iff minus-to-eq*
monoid.m-closed

monoid-axioms partial-object.select-convs(1) prod.collapse r-neg subset
subset-iff)

have $\text{snd } z \otimes (t' \otimes (t \otimes (\text{snd } x \otimes \text{fst } y))) = \text{snd } x \otimes (t \otimes (t' \otimes (\text{snd } z \otimes \text{fst } y)))$
using *comm-monoid-axioms-def[of R] f1 f3 subset a3 a4 a5 m-assoc*
by (*smt m-lcomm mem-Sigma-iff partial-object.select-convs(1) partial-object-ext-def*
rel-def

semiring-simprules(3) rev-subsetD surjective-pairing)

then have $\text{snd } z \otimes (t' \otimes (t \otimes (\text{snd } y \otimes \text{fst } x))) \ominus \text{snd } z \otimes (t' \otimes (t \otimes (\text{snd } x \otimes \text{fst } y))) \oplus$
 $\text{snd } x \otimes (t \otimes (t' \otimes (\text{snd } z \otimes \text{fst } y))) \ominus \text{snd } x \otimes (t \otimes (t' \otimes (\text{snd } y \otimes \text{fst } z)))$
 $=$
 $\text{snd } z \otimes (t' \otimes (t \otimes (\text{snd } y \otimes \text{fst } x))) \ominus \text{snd } x \otimes (t \otimes (t' \otimes (\text{snd } y \otimes \text{fst } z)))$
using *add.l-inv*
by (*smt a3 a4 a5 f1 f3 f5 is-abelian-group local.semiring-axioms mem-Sigma-iff*
minus-to-eq

monoid.m-closed monoid-axioms partial-object.select-convs(1) prod.collapse
rel-def

semiring.semiring-simprules(6) subset subset-iff)

then have $f7:\text{snd } z \otimes (t' \otimes (t \otimes (\text{snd } y \otimes \text{fst } x))) \ominus \text{snd } x \otimes (t \otimes (t' \otimes (\text{snd } y \otimes \text{fst } z))) = \mathbf{0}$
using $f5\ f6$
by (*smt* $\langle \text{snd } z \otimes (t' \otimes (t \otimes (\text{snd } x \otimes \text{fst } y))) = \text{snd } x \otimes (t \otimes (t' \otimes (\text{snd } z \otimes \text{fst } y))) \rangle$
 $\langle t' \otimes (\text{snd } z \otimes \text{fst } y) \ominus t' \otimes (\text{snd } y \otimes \text{fst } z) = \mathbf{0} \rangle$ *a4 a5 f3 is-abelian-group*
mem-Sigma-iff

minus-to-eq partial-object.select-convs(1) prod.collapse rel-def semir-
ing-simprules(3)
subset subset-iff)

moreover have $(t \otimes t' \otimes \text{snd } y) \otimes ((\text{snd } z \otimes \text{fst } x) \ominus (\text{snd } x \otimes \text{fst } z)) = ((t \otimes t' \otimes \text{snd } y) \otimes (\text{snd } z \otimes \text{fst } x)) \ominus ((t \otimes t' \otimes \text{snd } y) \otimes (\text{snd } x \otimes \text{fst } z))$

```

using r-distr f1 f3 subset a3 a4 a5 rel-def a-minus-def r-minus
by (smt SigmaE abelian-group.a-inv-closed is-abelian-group monoid.m-closed
monoid-axioms
  partial-object.select-convs(1) prod.sel(1) prod.sel(2) subset-iff)
moreover have f8:(t ⊗ t' ⊗ snd y) ⊗ (snd z ⊗ fst x) = snd z ⊗ (t' ⊗ (t ⊗
(snd y ⊗ fst x)))
using m-assoc comm-monoid-axioms-def[of R] f1 f3 subset a3 a4 a5 rel-def
rev-subsetD
by (smt SigmaE local.semiring-axioms m-lcomm partial-object.select-convs(1)
prod.sel(1)
  prod.sel(2) semiring.semiring-simprules(3))
moreover have f9:(t ⊗ t' ⊗ snd y) ⊗ (snd x ⊗ fst z) = snd x ⊗ (t ⊗ (t' ⊗
(snd y ⊗ fst z)))
using m-assoc comm-monoid-axioms-def[of R] f1 f3 subset a3 a4 a5 rel-def
rev-subsetD
by (smt SigmaE m-comm monoid.m-closed monoid-axioms partial-object.select-convs(1)
prod.sel(1)
  prod.sel(2))
then have f10:(t ⊗ t' ⊗ snd y) ⊗ (snd z ⊗ fst x) ⊖ (t ⊗ t' ⊗ snd y) ⊗ (snd
x ⊗ fst z) = 0
using f7 f8 f9
by simp
moreover have t ⊗ t' ⊗ snd y ∈ S
using f1 f3 a4 rel-def m-closed
by (simp add: mem-Times-iff)
then have (t ⊗ t' ⊗ snd y) ⊗ (snd z ⊗ fst x ⊖ snd x ⊗ fst z) = 0
using r-distr subset rev-subsetD f10 calculation(2)
by auto
thus x .=rel z
using rel-def ⟨t ⊗ t' ⊗ snd y ∈ S⟩
by auto
qed
qed

```

definition *eq-class-of-rng-of-frac*:: - ⇒ 'a ⇒ 'b ⇒ -set (infix <|₁> 10)
where $r \mid_{rel} s \equiv \{(r', s') \in \text{carrier } rel. (r, s) \text{ .=}_{rel} (r', s')\}$

lemma *class-of-to-rel*:
shows $\text{class-of}_{rel} (r, s) = (r \mid_{rel} s)$
using *eq-class-of-def*[of rel] *eq-class-of-rng-of-frac-def*[of rel]
by auto

lemma (in *eq-obj-rng-of-frac*) *zero-in-mult-submonoid*:
assumes $0 \in S$ **and** $(r, s) \in \text{carrier } rel$ **and** $(r', s') \in \text{carrier } rel$
shows $(r \mid_{rel} s) = (r' \mid_{rel} s')$
proof
show $(r \mid_{rel} s) \subseteq (r' \mid_{rel} s')$
proof
fix x

```

assume a1:  $x \in (r \mid_{rel} s)$ 
have  $\mathbf{0} \otimes (s' \otimes fst\ x \ominus snd\ x \otimes r') = \mathbf{0}$ 
  using l-zero subset rel-def a1 eq-class-of-rng-of-frac-def
  by (smt abelian-group.minus-closed assms(3) is-abelian-group l-null mem-Collect-eq
mem-Sigma-iff
  monoid.m-closed monoid-axioms old.prod.case partial-object.select-convs(1)
subset-iff surjective-pairing)
  thus  $x \in (r' \mid_{rel} s')$ 
  using assms(1) assms(3) rel-def eq-class-of-rng-of-frac-def
  by (smt SigmaE a1 eq-object.select-convs(1) l-null mem-Collect-eq minus-closed
old.prod.case
  partial-object.select-convs(1) prod.collapse semiring-simprules(3) subset
subset-iff)
qed
show  $(r' \mid_{rel} s') \subseteq (r \mid_{rel} s)$ 
proof
  fix x
  assume a1:  $x \in (r' \mid_{rel} s')$ 
  have  $\mathbf{0} \otimes (s \otimes fst\ x \ominus snd\ x \otimes r) = \mathbf{0}$ 
  using l-zero subset rel-def a1 eq-class-of-rng-of-frac-def
  by (metis (no-types, lifting) BNF-Def.Collect-case-prodD assms(2) l-null
mem-Sigma-iff
  minus-closed partial-object.select-convs(1) semiring-simprules(3) rev-subsetD)
  thus  $x \in (r \mid_{rel} s)$ 
  using assms(1) assms(2) rel-def eq-class-of-rng-of-frac-def
  by (smt SigmaE a1 eq-object.select-convs(1) l-null mem-Collect-eq minus-closed
old.prod.case
  partial-object.select-convs(1) prod.collapse semiring-simprules(3) subset
subset-iff)
qed
qed

```

definition *set-eq-class-of-rng-of-frac*:: $- \Rightarrow$ $-set$ ($\langle set'-class'-of \rangle$)
where $set-class-of_{rel} \equiv \{(r \mid_{rel} s) \mid r\ s, (r, s) \in carrier\ rel\}$

lemma *elem-eq-class*:

```

assumes equivalence S and  $x \in carrier\ S$  and  $y \in carrier\ S$  and  $x .=_S y$ 
shows  $class-of_S\ x = class-of_S\ y$ 
proof
show  $class-of_S\ x \subseteq class-of_S\ y$ 
proof
  fix z
  assume  $z \in class-of_S\ x$ 
  then have  $y .=_S z$ 
  using assms eq-class-of-def[of S x] equivalence.sym[of S x y] equivalence.trans
  by (metis (mono-tags, lifting) mem-Collect-eq)
thus  $z \in class-of_S\ y$ 
  using  $\langle z \in class-of_S\ x \rangle$ 

```

```

    by (simp add: eq-class-of-def)
qed
show class-ofS y ⊆ class-ofS x
proof
  fix z
  assume z ∈ class-ofS y
  then have x .=S z
    using assms eq-class-of-def equivalence.trans
    by (metis (mono-tags, lifting) mem-Collect-eq)
  thus z ∈ class-ofS x
    using ⟨z ∈ class-ofS y⟩
    by (simp add: eq-class-of-def)
qed
qed

lemma (in abelian-group) four-elem-comm:
  assumes a ∈ carrier G and b ∈ carrier G and c ∈ carrier G and d ∈ carrier
  G
  shows a ⊖ c ⊕ b ⊖ d = a ⊕ b ⊖ c ⊖ d
  using assms a-assoc a-comm
  by (simp add: a-minus-def)

lemma (in abelian-monoid) right-add-eq:
  assumes a = b
  shows c ⊕ a = c ⊕ b
  using assms
  by simp

lemma (in abelian-monoid) right-minus-eq:
  assumes a = b
  shows c ⊖ a = c ⊖ b
  by (simp add: assms)

lemma (in abelian-group) inv-add:
  assumes a ∈ carrier G and b ∈ carrier G
  shows ⊖ (a ⊕ b) = ⊖ a ⊖ b
  using assms minus-add
  by (simp add: a-minus-def)

lemma (in abelian-group) right-inv-add:
  assumes a ∈ carrier G and b ∈ carrier G and c ∈ carrier G
  shows c ⊖ a ⊖ b = c ⊖ (a ⊕ b)
  using assms
  by (simp add: a-minus-def add.m-assoc local.minus-add)

context eq-obj-rng-of-frac
begin

definition carrier-rng-of-frac:: - partial-object

```

where $\text{carrier-rng-of-frac} \equiv (\text{carrier} = \text{set-class-of_rel})$

definition $\text{mult-rng-of-frac}:: [-\text{set}, -\text{set}] \Rightarrow -\text{set}$

where $\text{mult-rng-of-frac } X Y \equiv$

$\text{let } x' = (\text{SOME } x. x \in X) \text{ in}$

$\text{let } y' = (\text{SOME } y. y \in Y) \text{ in}$

$(\text{fst } x' \otimes \text{fst } y')|_{\text{rel}} (\text{snd } x' \otimes \text{snd } y')$

definition $\text{rec-monoid-rng-of-frac}:: - \text{ monoid}$

where $\text{rec-monoid-rng-of-frac} \equiv (\text{carrier} = \text{set-class-of_rel}, \text{mult} = \text{mult-rng-of-frac},$
 $\text{one} = (\mathbf{1}|_{\text{rel}} \mathbf{1}))$

lemma $\text{member-class-to-carrier}$:

assumes $x \in (r |_{\text{rel}} s)$ **and** $y \in (r' |_{\text{rel}} s')$

shows $(\text{fst } x \otimes \text{fst } y, \text{snd } x \otimes \text{snd } y) \in \text{carrier } \text{rel}$

using $\text{assms } \text{rel-def } \text{eq-class-of-rng-of-frac-def}$

by $(\text{metis } (\text{no-types}, \text{lifting}) \text{Product-Type.Collect-case-prodD } m\text{-closed } \text{mem-Sigma-iff}$

$\text{partial-object.select-convs}(1) \text{ semiring-simprules}(3))$

lemma $\text{member-class-to-member-class}$:

assumes $x \in (r |_{\text{rel}} s)$ **and** $y \in (r' |_{\text{rel}} s')$

shows $(\text{fst } x \otimes \text{fst } y |_{\text{rel}} \text{snd } x \otimes \text{snd } y) \in \text{set-class-of_rel}$

using $\text{assms } \text{member-class-to-carrier}[\text{of } x \ r \ s \ y \ r' \ s'] \text{ set-eq-class-of-rng-of-frac-def}[\text{of}$
 $\text{rel}]$

$\text{eq-class-of-rng-of-frac-def}$

by auto

lemma $\text{closed-mult-rng-of-frac}$:

assumes $(r, s) \in \text{carrier } \text{rel}$ **and** $(t, u) \in \text{carrier } \text{rel}$

shows $(r |_{\text{rel}} s) \otimes_{\text{rec-monoid-rng-of-frac}} (t |_{\text{rel}} u) \in \text{set-class-of_rel}$

proof –

have $(r, s) \text{.}=_{\text{rel}} (r, s)$

using $\text{assms}(1) \text{equiv-obj-rng-of-frac equivalence-def}[\text{of } \text{rel}]$

by blast

then have $(r, s) \in (r |_{\text{rel}} s)$

using $\text{assms}(1)$

by $(\text{simp add: eq-class-of-rng-of-frac-def})$

then have $f1:\exists x. x \in (r |_{\text{rel}} s)$

by auto

have $f2:\exists y. y \in (t |_{\text{rel}} u)$

using $\text{assms}(2) \text{equiv-obj-rng-of-frac equivalence.refl eq-class-of-rng-of-frac-def}$

by fastforce

show $(r |_{\text{rel}} s) \otimes_{\text{rec-monoid-rng-of-frac}} (t |_{\text{rel}} u) \in \text{set-class-of_rel}$

using $f1 \ f2 \ \text{rec-monoid-rng-of-frac-def } \text{mult-rng-of-frac-def}[\text{of } (r |_{\text{rel}} s) (t |_{\text{rel}}$
 $u)]$

$\text{set-eq-class-of-rng-of-frac-def}[\text{of } \text{rel}] \ \text{member-class-to-member-class}[\text{of } x' \ r \ s \ y'$
 $t \ u]$

by $(\text{metis } (\text{mono-tags}, \text{lifting}) \text{mem-Collect-eq } \text{member-class-to-carrier } \text{monoid.select-convs}(1))$

someI-ex)

qed

lemma *non-empty-class*:
assumes $(r, s) \in \text{carrier } rel$
shows $(r \mid_{rel} s) \neq \{\}$
using *assms eq-class-of-rng-of-frac-def equiv-obj-rng-of-frac equivalence.refl*
by *fastforce*

lemma *mult-rng-of-frac-fundamental-lemma*:
assumes $(r, s) \in \text{carrier } rel$ **and** $(r', s') \in \text{carrier } rel$
shows $(r \mid_{rel} s) \otimes_{\text{rec-monoid-rng-of-frac}} (r' \mid_{rel} s') = (r \otimes r' \mid_{rel} s \otimes s')$
proof –
have $f1:(r \mid_{rel} s) \neq \{\}$
using *assms(1) non-empty-class*
by *auto*
have $(r' \mid_{rel} s') \neq \{\}$
using *assms(2) non-empty-class*
by *auto*
then have $\exists x \in (r \mid_{rel} s). \exists x' \in (r' \mid_{rel} s'). (r \mid_{rel} s) \otimes_{\text{rec-monoid-rng-of-frac}} (r' \mid_{rel} s') =$
 $(fst\ x \otimes fst\ x' \mid_{rel} snd\ x \otimes snd\ x')$
using *f1 rec-monoid-rng-of-frac-def*
by *(metis monoid.select-convs(1) mult-rng-of-frac-def some-in-eq)*
then obtain x **and** x' **where** $f2:x \in (r \mid_{rel} s)$ **and** $f3:x' \in (r' \mid_{rel} s')$
and $(r \mid_{rel} s) \otimes_{\text{rec-monoid-rng-of-frac}} (r' \mid_{rel} s') = (fst\ x \otimes fst\ x' \mid_{rel} snd\ x \otimes$
 $snd\ x')$
by *blast*
then have $(r, s) \text{.}=_rel (fst\ x, snd\ x)$
using *rel-def*
by *(metis (no-types, lifting) Product-Type.Collect-case-prodD eq-class-of-rng-of-frac-def)*
then obtain t **where** $f4:t \in S$ **and** $f5:t \otimes ((snd\ x \otimes r) \ominus (s \otimes fst\ x)) = \mathbf{0}$
using *rel-def*
by *auto*
have $(r', s') \text{.}=_rel (fst\ x', snd\ x')$
using *rel-def f3*
by *(metis (no-types, lifting) Product-Type.Collect-case-prodD eq-class-of-rng-of-frac-def)*
then obtain t' **where** $f6:t' \in S$ **and** $f7:t' \otimes (snd\ x' \otimes r' \ominus s' \otimes fst\ x') = \mathbf{0}$
using *rel-def*
by *auto*
have $f8:t \in \text{carrier } R$
using *f4 subset rev-subsetD*
by *auto*
have $f9:snd\ x \otimes r \in \text{carrier } R$
using *subset rev-subsetD f2 assms(1)*
by *(metis (no-types, lifting) BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def*
mem-Sigma-iff
partial-object.select-convs(1) rel-def semiring-simprules(3))

have $f10: \ominus (s \otimes \text{fst } x) \in \text{carrier } R$
using $\text{assms}(1)$ $\text{subset rev-subsetD } f2$
by ($\text{metis (no-types, lifting) BNF-Def.Collect-case-prodD abelian-group.a-inv-closed}$

$\text{eq-class-of-rng-of-frac-def is-abelian-group mem-Sigma-iff monoid.m-closed}$
 monoid-axioms
 $\text{partial-object.select-convs}(1)$ rel-def)
then have $t \otimes (\text{snd } x \otimes r) \ominus t \otimes (s \otimes \text{fst } x) = \mathbf{0}$
using $f8$ $f9$ $f10$ $f5$ $r\text{-distr}[of \text{snd } x \otimes r \ominus (s \otimes \text{fst } x) t]$ $a\text{-minus-def}$ $r\text{-minus}[of t s \otimes \text{fst } x]$
by ($\text{smt BNF-Def.Collect-case-prodD assms}(1)$ $\text{eq-class-of-rng-of-frac-def } f2$
 mem-Sigma-iff
 $\text{partial-object.select-convs}(1)$ $\text{rel-def semiring-simprules}(3)$ subset subset-iff)
then have $f11: t \otimes (\text{snd } x \otimes r) = t \otimes (s \otimes \text{fst } x)$
by ($\text{smt BNF-Def.Collect-case-prodD assms}(1)$ $\text{eq-class-of-rng-of-frac-def } f2$ $f8$
 is-abelian-group
 $\text{mem-Sigma-iff minus-to-eq monoid.m-closed monoid-axioms partial-object.select-convs}(1)$
 $\text{rel-def subset subset-iff}$)
have $f12: t' \in \text{carrier } R$
using $f6$ $\text{subset rev-subsetD}$
by auto
have $f13: \text{snd } x' \otimes r' \in \text{carrier } R$
using $\text{assms}(2)$ $f3$ $\text{subset rev-subsetD}$
by ($\text{metis (no-types, lifting) Product-Type.Collect-case-prodD eq-class-of-rng-of-frac-def}$

$\text{mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs}(1)$
 rel-def)
have $f14: \ominus (s' \otimes \text{fst } x') \in \text{carrier } R$
using $\text{assms}(2)$ $f3$ $\text{subset rev-subsetD}$
by ($\text{metis (no-types, lifting) BNF-Def.Collect-case-prodD abelian-group.a-inv-closed}$

$\text{eq-class-of-rng-of-frac-def is-abelian-group mem-Sigma-iff monoid.m-closed}$
 monoid-axioms
 $\text{partial-object.select-convs}(1)$ rel-def)
then have $t' \otimes (\text{snd } x' \otimes r') \ominus t' \otimes (s' \otimes \text{fst } x') = \mathbf{0}$
using $f12$ $f13$ $f14$ $f7$ $r\text{-distr}[of \text{snd } x' \otimes r' \ominus (s' \otimes \text{fst } x') t']$ $a\text{-minus-def}$
 $r\text{-minus}[of t' s' \otimes \text{fst } x']$
by ($\text{smt BNF-Def.Collect-case-prodD assms}(2)$ $\text{eq-class-of-rng-of-frac-def } f3$
 mem-Sigma-iff
 $\text{partial-object.select-convs}(1)$ $\text{rel-def semiring-simprules}(3)$ subset subset-iff)
then have $f15: t' \otimes (\text{snd } x' \otimes r') = t' \otimes (s' \otimes \text{fst } x')$
by ($\text{smt BNF-Def.Collect-case-prodD assms}(2)$ $\text{eq-class-of-rng-of-frac-def } f3$ $f12$
 is-abelian-group
 $\text{mem-Sigma-iff minus-to-eq monoid.m-closed monoid-axioms partial-object.select-convs}(1)$
 $\text{rel-def subset subset-iff}$)
have $t' \otimes t \in S$
using $f4$ $f6$ $m\text{-closed}$
by auto
then have $f16: t' \otimes t \in \text{carrier } R$

```

using subset rev-subsetD
by auto
have f17:(snd x ⊗ snd x') ⊗ (r ⊗ r') ∈ carrier R
using assms f2 f3
by (metis (no-types, lifting) BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def
mem-Sigma-iff
monoid.m-closed monoid-axioms partial-object.select-convs(1) rel-def subset
subset-iff)
have f18:(s ⊗ s') ⊗ (fst x ⊗ fst x') ∈ carrier R
using assms f2 f3
by (metis (no-types, lifting) BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def
mem-Sigma-iff
monoid.m-closed monoid-axioms partial-object.select-convs(1) rel-def subset
subset-iff)
then have f19:(t' ⊗ t) ⊗ ((snd x ⊗ snd x') ⊗ (r ⊗ r') ⊖ (s ⊗ s') ⊗ (fst x ⊗
fst x')) =
((t' ⊗ t) ⊗ (snd x ⊗ snd x')) ⊗ (r ⊗ r') ⊖ (t' ⊗ t) ⊗ ((s ⊗ s') ⊗ (fst x ⊗ fst
x'))
using f16 f17 f18 r-distr m-assoc r-minus a-minus-def
by (smt BNF-Def.Collect-case-prodD assms(1) assms(2) eq-class-of-rng-of-frac-def
f14 f2 f3
m-comm mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1)
rel-def
subset subset-iff)
then have f20:(t' ⊗ t) ⊗ (snd x ⊗ snd x') ⊗ (r ⊗ r') = (t' ⊗ t) ⊗ (snd x ⊗ r
⊗ snd x' ⊗ r')
using m-assoc m-comm f16 assms rel-def f2 f3
by (smt BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def mem-Sigma-iff

partial-object.select-convs(1) semiring-simprules(3) subset subset-iff)
then have ((t' ⊗ t) ⊗ (snd x ⊗ snd x')) ⊗ (r ⊗ r') = t' ⊗ ((t ⊗ snd x ⊗ r) ⊗
snd x' ⊗ r')
using m-assoc assms f2 f3 rel-def f8 f12
by (smt BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def mem-Sigma-iff
monoid.m-closed
monoid-axioms partial-object.select-convs(1) subset subset-iff)
then have f21:((t' ⊗ t) ⊗ (snd x ⊗ snd x')) ⊗ (r ⊗ r') = t' ⊗ (t ⊗ s ⊗ fst x)
⊗ snd x' ⊗ r'
using f11 m-assoc
by (smt BNF-Def.Collect-case-prodD assms(1) assms(2) eq-class-of-rng-of-frac-def
f12 f2 f3 f8
mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1)
rel-def subset subset-iff)
moreover have (t' ⊗ t) ⊗ ((s ⊗ s') ⊗ (fst x ⊗ fst x')) = (t' ⊗ s' ⊗ fst x') ⊗ t
⊗ s ⊗ fst x
using assms f2 f3 f8 f12 m-assoc m-comm rel-def
by (smt BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def mem-Sigma-iff
monoid.m-closed
monoid-axioms partial-object.select-convs(1) subset subset-iff)

```

then have $(t' \otimes t) \otimes ((s \otimes s') \otimes (fst\ x \otimes fst\ x')) = (t' \otimes snd\ x' \otimes r') \otimes t \otimes s$
 $\otimes\ fst\ x$
using *f15 m-assoc*
by (*smt BNF-Def.Collect-case-prodD assms(2) eq-class-of-rng-of-frac-def f12*
f3 mem-Sigma-iff
partial-object.select-convs(1) rel-def subset subset-iff)
then have $f22:(t' \otimes t) \otimes ((s \otimes s') \otimes (fst\ x \otimes fst\ x')) = t' \otimes ((t \otimes snd\ x \otimes r)$
 $\otimes\ snd\ x' \otimes r')$
using *m-assoc m-comm assms*
by (*smt BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def f12 f2 f21 f3 f8*
mem-Sigma-iff
partial-object.select-convs(1) rel-def semiring-simprules(3) subset subset-iff)
then have $f23:(t' \otimes t) \otimes ((snd\ x \otimes snd\ x') \otimes (r \otimes r') \ominus (s \otimes s') \otimes (fst\ x \otimes$
 $fst\ x')) = 0$
using *f19 f21 f22*
by (*metis* $\langle t' \otimes t \otimes (snd\ x \otimes snd\ x') \otimes (r \otimes r') = t' \otimes (t \otimes snd\ x \otimes r \otimes snd$
 $x' \otimes r') \rangle$
a-minus-def f16 f18 r-neg semiring-simprules(3))
have $f24:(r \otimes r', s \otimes s') \in carrier\ rel$
using *assms rel-def*
by *auto*
have $f25:(fst\ x \otimes fst\ x', snd\ x \otimes snd\ x') \in carrier\ rel$
using *f2 f3 member-class-to-carrier*
by *auto*
then have $(r \otimes r', s \otimes s') \dot{=}_{rel} (fst\ x \otimes fst\ x', snd\ x \otimes snd\ x')$
using *f23 f24 rel-def* $\langle t' \otimes t \in S \rangle$
by *auto*
then have $class_of_{rel}\ (r \otimes r', s \otimes s') = class_of_{rel}\ (fst\ x \otimes fst\ x', snd\ x \otimes snd$
 $x')$
using *f24 f25 equiv-obj-rng-of-frac elem-eq-class[of rel (r \otimes r', s \otimes s') (fst\ x \otimes*
 $fst\ x', snd\ x \otimes snd\ x')]$
eq-class-of-rng-of-frac-def
by *auto*
then have $(r \otimes r' \mid_{rel} s \otimes s') = (fst\ x \otimes fst\ x' \mid_{rel} snd\ x \otimes snd\ x')$
using *class-of-to-rel[of rel]*
by *auto*
thus *?thesis*
using $\langle (r \mid_{rel} s) \otimes_{rec-monoid-rng-of-frac} (r' \mid_{rel} s') = (fst\ x \otimes fst\ x' \mid_{rel} snd\ x$
 $\otimes\ snd\ x') \rangle$
trans sym
by *auto*
qed

lemma *member-class-to-assoc:*

assumes $x \in (r \mid_{rel} s)$ **and** $y \in (t \mid_{rel} u)$ **and** $z \in (v \mid_{rel} w)$
shows $((fst\ x \otimes fst\ y) \otimes fst\ z \mid_{rel} (snd\ x \otimes snd\ y) \otimes snd\ z) = (fst\ x \otimes (fst\ y \otimes$
 $fst\ z) \mid_{rel} snd\ x \otimes (snd\ y \otimes snd\ z))$
using *assms m-assoc subset rel-def rev-subsetD*
by (*smt BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def mem-Sigma-iff*)

partial-object.select-convs(1)

lemma *assoc-mult-rng-of-frac:*

assumes $(r, s) \in \text{carrier rel}$ **and** $(t, u) \in \text{carrier rel}$ **and** $(v, w) \in \text{carrier rel}$
shows $((r \mid_{\text{rel}} s) \otimes_{\text{rec-monoid-rng-of-frac}} (t \mid_{\text{rel}} u)) \otimes_{\text{rec-monoid-rng-of-frac}} (v \mid_{\text{rel}} w) =$
 $(r \mid_{\text{rel}} s) \otimes_{\text{rec-monoid-rng-of-frac}} ((t \mid_{\text{rel}} u) \otimes_{\text{rec-monoid-rng-of-frac}} (v \mid_{\text{rel}} w))$

proof –

have $((r \otimes t) \otimes v, (s \otimes u) \otimes w) = (r \otimes (t \otimes v), s \otimes (u \otimes w))$
using *assms m-assoc*
by (*metis (no-types, lifting) mem-Sigma-iff partial-object.select-convs(1) rel-def rev-subsetD subset*)
then have $f1:((r \otimes t) \otimes v \mid_{\text{rel}} (s \otimes u) \otimes w) = (r \otimes (t \otimes v) \mid_{\text{rel}} s \otimes (u \otimes w))$
by *simp*
have $f2:((r \mid_{\text{rel}} s) \otimes_{\text{rec-monoid-rng-of-frac}} (t \mid_{\text{rel}} u)) \otimes_{\text{rec-monoid-rng-of-frac}} (v \mid_{\text{rel}} w) =$
 $((r \otimes t) \otimes v \mid_{\text{rel}} (s \otimes u) \otimes w)$
using *assms mult-rng-of-frac-fundamental-lemma rel-def*
by *auto*
have $f3:(r \mid_{\text{rel}} s) \otimes_{\text{rec-monoid-rng-of-frac}} ((t \mid_{\text{rel}} u) \otimes_{\text{rec-monoid-rng-of-frac}} (v \mid_{\text{rel}} w)) =$
 $(r \otimes (t \otimes v) \mid_{\text{rel}} s \otimes (u \otimes w))$
using *assms mult-rng-of-frac-fundamental-lemma rel-def*
by *auto*
thus *?thesis*
using $f1 f2 f3$
by *simp*

qed

lemma *left-unit-mult-rng-of-frac:*

assumes $(r, s) \in \text{carrier rel}$
shows $\mathbf{1}_{\text{rec-monoid-rng-of-frac}} \otimes_{\text{rec-monoid-rng-of-frac}} (r \mid_{\text{rel}} s) = (r \mid_{\text{rel}} s)$
using *assms subset rev-subsetD rec-monoid-rng-of-frac-def mult-rng-of-frac-fundamental-lemma[of 1 1 r s]*
l-one[of r] l-one[of s] rel-def
by *auto*

lemma *right-unit-mult-rng-of-frac:*

assumes $(r, s) \in \text{carrier rel}$
shows $(r \mid_{\text{rel}} s) \otimes_{\text{rec-monoid-rng-of-frac}} \mathbf{1}_{\text{rec-monoid-rng-of-frac}} = (r \mid_{\text{rel}} s)$
using *assms subset rev-subsetD rec-monoid-rng-of-frac-def mult-rng-of-frac-fundamental-lemma[of r s 1 1]*
r-one[of r] r-one[of s] rel-def
by *auto*

lemma *monoid-rng-of-frac:*

shows *monoid (rec-monoid-rng-of-frac)*
proof

show $\bigwedge x y. x \in \text{carrier } \text{rec-monoid-rng-of-frac} \implies$
 $y \in \text{carrier } \text{rec-monoid-rng-of-frac} \implies x \otimes_{\text{rec-monoid-rng-of-frac}} y \in \text{carrier}$
 $\text{rec-monoid-rng-of-frac}$
using $\text{rec-monoid-rng-of-frac-def closed-mult-rng-of-frac}$
by $(\text{smt mem-Collect-eq partial-object.select-convs}(1) \text{ set-eq-class-of-rng-of-frac-def})$
show $\bigwedge x y z. x \in \text{carrier } \text{rec-monoid-rng-of-frac} \implies$
 $y \in \text{carrier } \text{rec-monoid-rng-of-frac} \implies$
 $z \in \text{carrier } \text{rec-monoid-rng-of-frac} \implies$
 $x \otimes_{\text{rec-monoid-rng-of-frac}} y \otimes_{\text{rec-monoid-rng-of-frac}} z =$
 $x \otimes_{\text{rec-monoid-rng-of-frac}} (y \otimes_{\text{rec-monoid-rng-of-frac}} z)$
using $\text{assoc-mult-rng-of-frac}$
by $(\text{smt mem-Collect-eq partial-object.select-convs}(1) \text{ rec-monoid-rng-of-frac-def}$
 $\text{ set-eq-class-of-rng-of-frac-def})$
show $\mathbf{1}_{\text{rec-monoid-rng-of-frac}} \in \text{carrier } \text{rec-monoid-rng-of-frac}$
using $\text{rec-monoid-rng-of-frac-def rel-def set-eq-class-of-rng-of-frac-def}$
by fastforce
show $\bigwedge x. x \in \text{carrier } \text{rec-monoid-rng-of-frac} \implies \mathbf{1}_{\text{rec-monoid-rng-of-frac}} \otimes_{\text{rec-monoid-rng-of-frac}}$
 $x = x$
using $\text{left-unit-mult-rng-of-frac}$
by $(\text{smt mem-Collect-eq partial-object.select-convs}(1) \text{ rec-monoid-rng-of-frac-def}$
 $\text{ set-eq-class-of-rng-of-frac-def})$
show $\bigwedge x. x \in \text{carrier } \text{rec-monoid-rng-of-frac} \implies x \otimes_{\text{rec-monoid-rng-of-frac}} \mathbf{1}_{\text{rec-monoid-rng-of-frac}}$
 $= x$
using $\text{right-unit-mult-rng-of-frac}$
by $(\text{smt mem-Collect-eq partial-object.select-convs}(1) \text{ rec-monoid-rng-of-frac-def}$
 $\text{ set-eq-class-of-rng-of-frac-def})$
qed

lemma $\text{comm-mult-rng-of-frac}$:

assumes $(r, s) \in \text{carrier } \text{rel}$ **and** $(r', s') \in \text{carrier } \text{rel}$
shows $(r \mid_{\text{rel}} s) \otimes_{\text{rec-monoid-rng-of-frac}} (r' \mid_{\text{rel}} s') = (r' \mid_{\text{rel}} s') \otimes_{\text{rec-monoid-rng-of-frac}}$
 $(r \mid_{\text{rel}} s)$

proof –

have $f1: (r \mid_{\text{rel}} s) \otimes_{\text{rec-monoid-rng-of-frac}} (r' \mid_{\text{rel}} s') = (r \otimes r' \mid_{\text{rel}} s \otimes s')$

using $\text{assms mult-rng-of-frac-fundamental-lemma}$

by simp

have $f2: (r' \mid_{\text{rel}} s') \otimes_{\text{rec-monoid-rng-of-frac}} (r \mid_{\text{rel}} s) = (r' \otimes r \mid_{\text{rel}} s' \otimes s)$

using $\text{assms mult-rng-of-frac-fundamental-lemma}$

by simp

have $f3: r \otimes r' = r' \otimes r$

using $\text{assms rel-def m-comm}$

by simp

have $f4: s \otimes s' = s' \otimes s$

using $\text{assms rel-def subset rev-subsetD m-comm}$

by $(\text{metis (no-types, lifting) mem-Sigma-iff partial-object.select-convs}(1))$

thus $?thesis$

using $f1 f2 f3 f4$

by simp

qed

lemma *comm-monoid-rng-of-frac*:

shows *comm-monoid* (*rec-monoid-rng-of-frac*)

using *comm-monoid-def* *Group.comm-monoid-axioms-def* *monoid-rng-of-frac* *comm-mult-rng-of-frac*

by (*smt mem-Collect-eq partial-object.select-convs*(1) *rec-monoid-rng-of-frac-def*
set-eq-class-of-rng-of-frac-def)

definition *add-rng-of-frac*:: [-set, -set] \Rightarrow -set

where *add-rng-of-frac* *X Y* \equiv

let $x' = (\text{SOME } x. x \in X)$ *in*

let $y' = (\text{SOME } y. y \in Y)$ *in*

$(\text{snd } y' \otimes \text{fst } x' \oplus \text{snd } x' \otimes \text{fst } y') \mid_{\text{rel}} (\text{snd } x' \otimes \text{snd } y')$

definition *rec-rng-of-frac*:: - ring

where *rec-rng-of-frac* \equiv

$(\mid \text{carrier} = \text{set-class-of}_{\text{rel}}, \text{mult} = \text{mult-rng-of-frac}, \text{one} = (\mathbf{1} \mid_{\text{rel}} \mathbf{1}), \text{zero} = (\mathbf{0} \mid_{\text{rel}} \mathbf{1}), \text{add} = \text{add-rng-of-frac} \mid)$

lemma *add-rng-of-frac-fundamental-lemma*:

assumes $(r, s) \in \text{carrier } \text{rel}$ **and** $(r', s') \in \text{carrier } \text{rel}$

shows $(r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r' \mid_{\text{rel}} s') = (s' \otimes r \oplus s \otimes r' \mid_{\text{rel}} s \otimes s')$

proof –

have $\exists x' \in (r \mid_{\text{rel}} s). \exists y' \in (r' \mid_{\text{rel}} s'). (r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r' \mid_{\text{rel}} s') =$
 $(\text{snd } y' \otimes \text{fst } x' \oplus \text{snd } x' \otimes \text{fst } y' \mid_{\text{rel}} \text{snd } x' \otimes \text{snd } y')$

using *assms* *rec-rng-of-frac-def* *add-rng-of-frac-def*[*of* $(r \mid_{\text{rel}} s)$ $(r' \mid_{\text{rel}} s')$]

by (*metis non-empty-class ring-record-simps*(12) *some-in-eq*)

then obtain x' **and** y' **where** $f1: x' \in (r \mid_{\text{rel}} s)$ **and** $f2: y' \in (r' \mid_{\text{rel}} s')$ **and**

$f3: (r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r' \mid_{\text{rel}} s') = (\text{snd } y' \otimes \text{fst } x' \oplus \text{snd } x' \otimes \text{fst } y' \mid_{\text{rel}} \text{snd } x' \otimes \text{snd } y')$

by *auto*

then have $(r, s) \text{.}=\text{rel } x'$

using *f1* *rel-def* *eq-class-of-rng-of-frac-def*[*of* *rel* r s]

by *auto*

then obtain t **where** $f4: t \in S$ **and** $f5: t \otimes (\text{snd } x' \otimes r \ominus s \otimes \text{fst } x') = \mathbf{0}$

using *rel-def*

by *auto*

have $(r', s') \text{.}=\text{rel } y'$

using *f2* *rel-def* *eq-class-of-rng-of-frac-def*[*of* *rel* r' s']

by *auto*

then obtain t' **where** $f6: t' \in S$ **and** $f7: t' \otimes (\text{snd } y' \otimes r' \ominus s' \otimes \text{fst } y') = \mathbf{0}$

using *rel-def*

by *auto*

then have $f8: t \otimes t' \in S$

using *m-closed* *f4* *f6*

by *simp*

then have $(s' \otimes r \oplus s \otimes r', s \otimes s') \text{.}=\text{rel } (\text{snd } y' \otimes \text{fst } x' \oplus \text{snd } x' \otimes \text{fst } y', \text{snd } x' \otimes \text{snd } y')$

proof –

have $f9:t' \otimes s' \otimes \text{snd } y' \in \text{carrier } R$
using $f6 \text{ assms}(2) f2 \text{ subset rev-subsetD eq-class-of-rng-of-frac-def rel-def}$
by *fastforce*
have $f10:\text{snd } x' \otimes r \in \text{carrier } R$
using $\text{assms}(1) f1 \text{ rel-def subset rev-subsetD}$
by (*metis (no-types, lifting) Product-Type.Collect-case-prodD eq-class-of-rng-of-frac-def*)

mem-Sigma-iff partial-object.select-convs(1) semiring-simprules(3)
have $f11:s \otimes \text{fst } x' \in \text{carrier } R$
using $\text{assms}(1) \text{ subset rev-subsetD } f1 \text{ rel-def}$
by (*metis (no-types, lifting) Product-Type.Collect-case-prodD eq-class-of-rng-of-frac-def*)

mem-Sigma-iff partial-object.select-convs(1) semiring-simprules(3)
have $t \otimes (\text{snd } x' \otimes r \ominus s \otimes \text{fst } x') = t \otimes (\text{snd } x' \otimes r) \ominus t \otimes (s \otimes \text{fst } x')$
using $f9 f10 f11 f4 \text{ subset rev-subsetD } r\text{-distr}[of \text{snd } x' \otimes r s \otimes \text{fst } x' t]$
a-minus-def
r-minus[of t s \otimes fst x']
by (*smt add.inv-closed monoid.m-closed monoid-axioms r-distr*)
then have $f12:(t' \otimes s' \otimes \text{snd } y') \otimes (t \otimes (\text{snd } x' \otimes r \ominus s \otimes \text{fst } x')) =$
 $t' \otimes s' \otimes \text{snd } y' \otimes t \otimes (\text{snd } x' \otimes r) \ominus (t' \otimes s' \otimes \text{snd } y' \otimes t \otimes (s \otimes \text{fst } x'))$
using $f9 r\text{-distr}[of - - t' \otimes s' \otimes \text{snd } y'] \text{ rel-def } r\text{-minus } a\text{-minus-def}$
by (*smt abelian-group.minus-to-eq f10 f11 f4 f5 is-abelian-group m-assoc monoid.m-closed*)
monoid-axioms r-neg r-null subset subset-iff
have $f13:(\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r) \in \text{carrier } R$
using $\text{assms } f1 f2 \text{ subset rev-subsetD}$
by (*metis (no-types, lifting) BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def*)

mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1)
rel-def
have $f14:(s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \in \text{carrier } R$
using $\text{assms } f1 f2 \text{ subset rev-subsetD}$
by (*metis (no-types, lifting) BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def*)

mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1)
rel-def
then have $(t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r) \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x')) =$
 $(t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r)) \ominus (t \otimes t') \otimes ((s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x'))$
using $f13 f14 f8 \text{ subset rev-subsetD } r\text{-distr rel-def } r\text{-minus } a\text{-minus-def}$
by (*smt add.inv-closed semiring-simprules(3)*)
have $f15:s \otimes s' \in \text{carrier } R$
using $\text{assms rel-def subset rev-subsetD}$
by *auto*
have $f16:\text{snd } y' \otimes \text{fst } x' \in \text{carrier } R$
using $f1 f2 \text{ rel-def subset rev-subsetD}[of - S] \text{ monoid.m-closed}[of R \text{snd } y' \text{fst } x']$
by (*metis (no-types, lifting) BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def*)

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      mem-Sigma-iff monoid-axioms partial-object.select-convs(1))
have f17:t ∈ carrier R
  using f4 subset rev-subsetD
  by auto
have f18:t' ∈ carrier R
  using f6 subset rev-subsetD
  by auto
have f19:s ∈ carrier R
  using assms(1) rel-def subset
  by auto
have f20:s' ∈ carrier R
  using assms(2) rel-def subset
  by auto
have f21:snd y' ∈ carrier R
  using f2 rel-def subset rev-subsetD
by (metis (no-types, lifting) Product-Type.Collect-case-prodD eq-class-of-rng-of-frac-def

      mem-Sigma-iff partial-object.select-convs(1))
have f22:fst x' ∈ carrier R
  using f1 rel-def
  by (metis (no-types, lifting) Product-Type.Collect-case-prodD eq-class-of-rng-of-frac-def
mem-Sigma-iff
      partial-object.select-convs(1))
  then have f23:(t ⊗ t') ⊗ ((s ⊗ s') ⊗ (snd y' ⊗ fst x')) = t' ⊗ s' ⊗ snd y' ⊗
t ⊗ (s ⊗ fst x')
  using f17 f18 f19 f20 f21 m-assoc m-comm
  by (smt BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def f1 f4 f6 mem-Sigma-iff

      partial-object.select-convs(1) rel-def semiring-simprules(3) subset-iff)
  have f24:(t ⊗ t') ⊗ ((snd x' ⊗ snd y') ⊗ (s' ⊗ r)) = t' ⊗ s' ⊗ snd y' ⊗ t ⊗
(snd x' ⊗ r)
  using f17 f18 f20 f21 m-assoc m-comm
  by (smt BNF-Def.Collect-case-prodD assms(1) eq-class-of-rng-of-frac-def f1
f2 f4 f6
      mem-Sigma-iff partial-object.select-convs(1) rel-def semiring-simprules(3)
subset subset-iff)
  then have (t ⊗ t') ⊗ ((snd x' ⊗ snd y') ⊗ (s' ⊗ r)) ⊖ (t ⊗ t') ⊗ ((s ⊗ s') ⊗
(snd y' ⊗ fst x'))=
  (t' ⊗ s' ⊗ snd y' ⊗ t ⊗ (snd x' ⊗ r)) ⊖ (t' ⊗ s' ⊗ snd y' ⊗ t ⊗ (s ⊗ fst x'))
  using f23 f24
  by simp
  then have f25:(t' ⊗ s' ⊗ snd y') ⊗ (t ⊗ (snd x' ⊗ r ⊖ s ⊗ fst x')) =
  (t ⊗ t') ⊗ ((snd x' ⊗ snd y') ⊗ (s' ⊗ r)) ⊖ (t ⊗ t') ⊗ ((s ⊗ s') ⊗ (snd y' ⊗
fst x'))
  using f12
  by simp
  have f26:(t ⊗ t') ⊗ ((snd x' ⊗ snd y') ⊗ (s ⊗ r')) ⊖ (t ⊗ t') ⊗ ((s ⊗ s') ⊗
(snd x' ⊗ fst y')) =

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$t \otimes s \otimes \text{snd } x' \otimes t' \otimes (\text{snd } y' \otimes r') \ominus (t \otimes s \otimes \text{snd } x' \otimes t' \otimes (s' \otimes \text{fst } y'))$
by (*smt BNF-Def.Collect-case-prodD assms(2) eq-class-of-rng-of-frac-def f1 f17 f18 f19 f2*)
m-assoc m-comm mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1) rel-def subset subset-iff
have *f27*: $\text{snd } y' \otimes r' \in \text{carrier } R$
using *assms(2) f21 rel-def*
by *auto*
have *f28*: $s' \otimes \text{fst } y' \in \text{carrier } R$
using *f20 assms(2)*
by (*metis (no-types, lifting) BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def f2*)
mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1) rel-def
then have $t' \otimes (\text{snd } y' \otimes r' \ominus s' \otimes \text{fst } y') = t' \otimes (\text{snd } y' \otimes r') \ominus t' \otimes (s' \otimes \text{fst } y')$
using *f18 f27 f28 r-minus[of t' s' \otimes fst y']*
by (*simp add: a-minus-def r-distr*)
then have *f29*: $(t \otimes s \otimes \text{snd } x') \otimes (t' \otimes (\text{snd } y' \otimes r' \ominus s' \otimes \text{fst } y')) =$
 $(t \otimes s \otimes \text{snd } x') \otimes (t' \otimes (\text{snd } y' \otimes r') \ominus t' \otimes (s' \otimes \text{fst } y'))$
by *simp*
have $t \otimes s \otimes \text{snd } x' \in \text{carrier } R$
using *f17 f19 f1 subset assms(1) eq-class-of-rng-of-frac-def f4 rel-def*
by *fastforce*
then have *f30*: $(t \otimes s \otimes \text{snd } x') \otimes (t' \otimes (\text{snd } y' \otimes r' \ominus s' \otimes \text{fst } y')) =$
 $(t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r')) \ominus (t \otimes t') \otimes ((s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y'))$
using *f26 f29 r-distr*
by (*smt <t' \otimes (snd y' \otimes r' \ominus s' \otimes fst y') = t' \otimes (snd y' \otimes r') \ominus t' \otimes (s' \otimes fst y')>*)
a-minus-def abelian-group.minus-to-eq f18 f27 f28 f7 is-abelian-group m-assoc monoid.m-closed
monoid-axioms r-neg semiring-simprules(15))
then have *f31*: $((t' \otimes s' \otimes \text{snd } y') \otimes (t \otimes (\text{snd } x' \otimes r \ominus s \otimes \text{fst } x'))) \oplus ((t \otimes s \otimes \text{snd } x') \otimes (t' \otimes (\text{snd } y' \otimes r' \ominus s' \otimes \text{fst } y')))$
 $= ((t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r))) \ominus (t \otimes t') \otimes ((s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x')) \oplus$
 $((t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r')) \ominus (t \otimes t') \otimes ((s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y')))$
using *f25 f30*
by *simp*
have *f32*: $(t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r)) \ominus (t \otimes t') \otimes ((s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x'))$
 $= (t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r)) \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x')$
using *f17 f18 r-distr*
by (*simp add: <t \otimes t' \otimes (snd x' \otimes snd y' \otimes (s' \otimes r)) \ominus s \otimes s' \otimes (snd y' \otimes fst x') = t \otimes t' \otimes (snd x' \otimes snd y' \otimes (s' \otimes r)) \ominus t \otimes t' \otimes (s \otimes s' \otimes (snd y' \otimes fst x'))>*)
have *f33*: $(t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r')) \ominus (t \otimes t') \otimes ((s \otimes s') \otimes$

$(\text{snd } x' \otimes \text{fst } y') =$
 $(t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y'))$
using $r\text{-distr}[of - - t \otimes t']$ $f17$ $f18$ $a\text{-minus-def}$ $r\text{-minus}$
by $(\text{smt BNF-Def.Collect-case-prodD abelian-group.a-inv-closed assms}(1)$
 $\text{assms}(2)$
 $\text{eq-class-of-rng-of-frac-def } f1$ $f2$ $\text{is-abelian-group mem-Sigma-iff partial-object.select-convs}(1)$
 $\text{rel-def semiring-simprules}(3)$ $\text{subset subset-iff})$
have $f34:(\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r \oplus s \otimes r') = (\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes$
 $r) \oplus (\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r')$
using $r\text{-distr}$
by $(\text{metis (no-types, lifting) BNF-Def.Collect-case-prodD assms}(1)$ $\text{assms}(2)$
 $\text{eq-class-of-rng-of-frac-def}$
 $f1$ $f2$ $\text{mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs}(1)$
 rel-def
 $\text{subset subset-iff})$
then have $(t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r \oplus s \otimes r')) =$
 $(t \otimes t') \otimes (\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r) \oplus (t \otimes t') \otimes (\text{snd } x' \otimes \text{snd } y') \otimes (s$
 $\otimes r')$
by $(\text{smt BNF-Def.Collect-case-prodD assms}(1)$ $\text{assms}(2)$ $\text{eq-class-of-rng-of-frac-def}$
 $f1$ $f17$ $f18$
 $f2$ $m\text{-assoc mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs}(1)$
 $r\text{-distr rel-def subset subset-iff})$
have $f35:(s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x' \oplus \text{snd } x' \otimes \text{fst } y') = (s \otimes s') \otimes (\text{snd } y' \otimes$
 $\text{fst } x') \oplus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y')$
using $r\text{-distr } f19$ $f20$
by $(\text{metis (no-types, lifting) BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def}$
 $f1$ $f2$
 $\text{mem-Sigma-iff partial-object.select-convs}(1)$ $\text{rel-def semiring-simprules}(3)$
 $\text{subset subset-iff})$
then have $f36:(t \otimes t') \otimes (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x' \oplus \text{snd } x' \otimes \text{fst } y') =$
 $(t \otimes t') \otimes (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \oplus (t \otimes t') \otimes (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst}$
 $y')$
by $(\text{smt BNF-Def.Collect-case-prodD assms}(1)$ $\text{assms}(2)$ $\text{eq-class-of-rng-of-frac-def}$
 $f1$ $f17$ $f18$ $f2$
 $\text{mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs}(1)$
 $r\text{-distr rel-def}$
 $\text{subset subset-iff})$
have $f37:(t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r) \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst}$
 $x')) \in \text{carrier } R$
by $(\text{simp add: } f13$ $f14$ $f17$ $f18)$
have $f38:(t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst}$
 $y')) \in \text{carrier } R$
using $\langle t \otimes s \otimes \text{snd } x' \in \text{carrier } R \rangle$ $f30$ $f33$ $f7$ zero-closed
by auto
have $f39:(t \otimes t') \otimes ((\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r)) \ominus (t \otimes t') \otimes ((s \otimes s') \otimes$
 $(\text{snd } y' \otimes \text{fst } x')) \in \text{carrier } R$
by $(\text{simp add: } f32$ $f37)$
have $\text{snd } x' \otimes \text{snd } y' \in \text{carrier } R$

using *f1 f2 subset rev-subsetD*
by (*metis (no-types, lifting) BNF-Def.Collect-case-prodD eq-class-of-rng-of-frac-def*)

mem-Sigma-iff partial-object.select-convs(1) rel-def semiring-simprules(3)
have $(t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s' \otimes r) \ominus (s \otimes s') \otimes (snd\ y' \otimes fst\ x'))$
 \oplus

$(t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes (snd\ x' \otimes fst\ y')) =$
 $(t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s' \otimes r)) \ominus (t \otimes t') \otimes ((s \otimes s') \otimes (snd\ y' \otimes$
fst x')) \oplus
 $(t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s \otimes r')) \ominus (t \otimes t') \otimes ((s \otimes s') \otimes (snd\ x' \otimes$
fst y'))

using *f32 f33 ⟨snd x' ⊗ snd y' ∈ carrier R⟩ ⟨t ⊗ s ⊗ snd x' ∈ carrier R⟩*
assms(2) f17 f18 f19
f25 f30 f5 f7 f9 l-zero r-null rel-def zero-closed
apply *clarsimp*
using *l-zero semiring-simprules(3) by presburger*
then have *f40:((t' ⊗ s' ⊗ snd y') ⊗ (t ⊗ (snd x' ⊗ r ⊕ s ⊗ fst x'))) ⊕*
 $((t \otimes s \otimes snd\ x') \otimes (t' \otimes (snd\ y' \otimes r' \ominus s' \otimes fst\ y'))) =$
 $((t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s' \otimes r) \ominus (s \otimes s') \otimes (snd\ y' \otimes fst\ x'))) \oplus$
 $((t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes (snd\ x' \otimes fst\ y')))$
using *f31*
by (*simp add: f32 f33*)
have *f41:(snd x' ⊗ snd y') ⊗ (s' ⊗ r) ⊕ (s ⊗ s') ⊗ (snd y' ⊗ fst x') ∈ carrier*
R
by (*simp add: f13 f14*)
have *f42:(snd x' ⊗ snd y') ⊗ (s ⊗ r') ⊕ (s ⊗ s') ⊗ (snd x' ⊗ fst y') ∈ carrier*
R
by (*smt BNF-Def.Collect-case-prodD abelian-group.minus-closed assms(1)*)
assms(2)
eq-class-of-rng-of-frac-def f1 f2 is-abelian-group mem-Sigma-iff partial-object.select-convs(1)

rel-def semiring-simprules(3) subset subset-iff)
then have $(t' \otimes s' \otimes snd\ y') \otimes (t \otimes (snd\ x' \otimes r \oplus s \otimes fst\ x')) \oplus$
 $(t \otimes s \otimes snd\ x') \otimes (t' \otimes (snd\ y' \otimes r' \ominus s' \otimes fst\ y')) =$
 $(t \otimes t') \otimes (((snd\ x' \otimes snd\ y') \otimes (s' \otimes r) \ominus (s \otimes s') \otimes (snd\ y' \otimes fst\ x')) \oplus$
 $((snd\ x' \otimes snd\ y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes (snd\ x' \otimes fst\ y')))$
using *r-distr[of (snd x' ⊗ snd y') ⊗ (s' ⊗ r) ⊕ (s ⊗ s') ⊗ (snd y' ⊗ fst x')*
 $(snd\ x' \otimes snd\ y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes (snd\ x' \otimes fst\ y')\ t \otimes t']$
f17 f18 f40 f41 f42
by *simp*
have $(snd\ x' \otimes snd\ y') \otimes (s' \otimes r) \ominus (s \otimes s') \otimes (snd\ y' \otimes fst\ x') \oplus (snd\ x' \otimes$
 $snd\ y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes (snd\ x' \otimes fst\ y') =$
 $(snd\ x' \otimes snd\ y') \otimes (s' \otimes r) \oplus (snd\ x' \otimes snd\ y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes$
 $(snd\ y' \otimes fst\ x') \ominus (s \otimes s') \otimes (snd\ x' \otimes fst\ y')$
using *four-elem-comm[of (snd x' ⊗ snd y') ⊗ (s' ⊗ r) (snd x' ⊗ snd y') ⊗*
 $(s \otimes r') (s \otimes s') \otimes (snd\ y' \otimes fst\ x') (s \otimes s') \otimes (snd\ x' \otimes fst\ y')]$
by (*smt BNF-Def.Collect-case-prodD assms eq-class-of-rng-of-frac-def f1 f2*)
mem-Sigma-iff partial-object.select-convs(1) rel-def semiring-simprules(3)
subset subset-iff)

then have $(\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r) \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \oplus (\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y') =$
 $((\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r) \oplus (\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r')) \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y')$
by blast
then have $f43:(\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r) \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \oplus (\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y') =$
 $(\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r \oplus s \otimes r') \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y')$
using f34
by simp
have $(\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r') \in \text{carrier } R$
using $\langle \text{snd } x' \otimes \text{snd } y' \in \text{carrier } R \rangle \text{ assms}(2) f19 \text{ rel-def}$
by auto
have $(s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y') \in \text{carrier } R$
by $(\text{metis } (\text{no-types, lifting}) \text{BNF-Def.Collect-case-prodD assms eq-class-of-rng-of-frac-def f1 f2 mem-Sigma-iff partial-object.select-convs}(1) \text{ rel-def}$
 $\text{semiring-simprules}(3) \text{ subset subset-iff})$
then have $f43\text{bis}:(\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r) \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \oplus ((\text{snd } x' \otimes \text{snd } y') \otimes (s \otimes r') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y')) =$
 $(\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r \oplus s \otimes r') \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y')$
using a-assoc a-minus-def f41 f43
by $(\text{smt } \langle \text{snd } x' \otimes \text{snd } y' \otimes (s \otimes r') \in \text{carrier } R \rangle \text{ add.l-inv-ex add.m-closed minus-equality})$
have $f44:s \otimes s' \otimes (\text{snd } y' \otimes \text{fst } x') \in \text{carrier } R$
by (simp add: f14)
have $f45:s \otimes s' \otimes (\text{snd } x' \otimes \text{fst } y') \in \text{carrier } R$
by $(\text{metis } (\text{no-types, lifting}) \text{BNF-Def.Collect-case-prodD assms eq-class-of-rng-of-frac-def f1 f2 mem-Sigma-iff partial-object.select-convs}(1) \text{ rel-def}$
 $\text{semiring-simprules}(3) \text{ subset subset-iff})$
then have $\ominus ((s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \oplus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y')) =$
 $\ominus ((s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x')) \ominus ((s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y'))$
using f44 f45 inv-add
by auto
then have $\ominus ((s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \oplus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y')) =$
 $\ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y')$
using l-minus[of s \otimes s']
by $(\text{simp add: a-minus-def f15 f16 f45})$
then have $(\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r \oplus s \otimes r') \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y') =$
 $(\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r \oplus s \otimes r') \ominus ((s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \oplus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y'))$
using right-inv-add $\langle \text{snd } x' \otimes \text{snd } y' \in \text{carrier } R \rangle \text{ assms}(2) f13 f19 f34 f44 f45 \text{ rel-def}$
by auto
then have $(\text{snd } x' \otimes \text{snd } y') \otimes (s' \otimes r \oplus s \otimes r') \ominus (s \otimes s') \otimes (\text{snd } y' \otimes \text{fst } x') \ominus (s \otimes s') \otimes (\text{snd } x' \otimes \text{fst } y')$

$x' \oplus (s \otimes s') \otimes (snd\ x' \otimes fst\ y') =$
 $(snd\ x' \otimes snd\ y') \otimes (s' \otimes r \oplus s \otimes r') \oplus ((s \otimes s') \otimes (snd\ y' \otimes fst\ x' \oplus snd$
 $x' \otimes fst\ y'))$
using *r-distr*
by (*simp add: f35*)
then have $((snd\ x' \otimes snd\ y') \otimes (s' \otimes r) \oplus (s \otimes s') \otimes (snd\ y' \otimes fst\ x')) \oplus$
 $((snd\ x' \otimes snd\ y') \otimes (s \otimes r') \oplus (s \otimes s') \otimes (snd\ x' \otimes fst\ y'))$
 $= (snd\ x' \otimes snd\ y') \otimes (s' \otimes r \oplus s \otimes r') \oplus ((s \otimes s') \otimes (snd\ y' \otimes fst\ x' \oplus snd$
 $x' \otimes fst\ y'))$
using *f43bis*
by *simp*
then have $(t \otimes t') \otimes (((snd\ x' \otimes snd\ y') \otimes (s' \otimes r) \oplus (s \otimes s') \otimes (snd\ y' \otimes$
 $fst\ x')) \oplus ((snd\ x' \otimes snd\ y') \otimes (s \otimes r') \oplus (s \otimes s') \otimes (snd\ x' \otimes fst\ y')))$
 $= (t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s' \otimes r \oplus s \otimes r') \oplus ((s \otimes s') \otimes (snd\ y' \otimes$
 $fst\ x' \oplus snd\ x' \otimes fst\ y')))$
by *simp*
then have $(t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s' \otimes r) \oplus (s \otimes s') \otimes (snd\ y' \otimes fst$
 $x')) \oplus$
 $(t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s \otimes r') \oplus (s \otimes s') \otimes (snd\ x' \otimes fst\ y')) =$
 $(t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s' \otimes r \oplus s \otimes r') \oplus ((s \otimes s') \otimes (snd\ y' \otimes$
 $fst\ x' \oplus snd\ x' \otimes fst\ y')))$
using *r-distr[of - - t \otimes t'] f17 f18 <t' \otimes s' \otimes snd\ y' \otimes (t \otimes (snd\ x' \otimes r \oplus s
 $\otimes fst\ x')) \oplus t \otimes s \otimes snd\ x' \otimes (t' \otimes (snd\ y' \otimes r' \oplus s' \otimes fst\ y')) = t \otimes t' \otimes (snd$
 $x' \otimes snd\ y' \otimes (s' \otimes r) \oplus s \otimes s' \otimes (snd\ y' \otimes fst\ x') \oplus (snd\ x' \otimes snd\ y' \otimes (s \otimes r')$
 $\oplus s \otimes s' \otimes (snd\ x' \otimes fst\ y'))\rangle$* *f40*
by *auto*
then have $(t' \otimes s' \otimes snd\ y') \otimes (t \otimes (snd\ x' \otimes r \oplus s \otimes fst\ x')) \oplus$
 $(t \otimes s \otimes snd\ x') \otimes (t' \otimes (snd\ y' \otimes r' \oplus s' \otimes fst\ y')) =$
 $(t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s' \otimes r \oplus s \otimes r') \oplus (s \otimes s') \otimes (snd\ y' \otimes fst$
 $x' \oplus snd\ x' \otimes fst\ y'))$
using *f40*
by *simp*
then have $(t \otimes t') \otimes ((snd\ x' \otimes snd\ y') \otimes (s' \otimes r \oplus s \otimes r') \oplus (s \otimes s') \otimes$
 $(snd\ y' \otimes fst\ x' \oplus snd\ x' \otimes fst\ y')) = 0$
using *f5 f7*
by (*simp add: <t \otimes s \otimes snd\ x' \in carrier R> f9*)
thus *?thesis*
using *rel-def f8*
by *auto*
qed
then have $(s' \otimes r \oplus s \otimes r' \mid_{rel\ s \otimes s'}) = (snd\ y' \otimes fst\ x' \oplus snd\ x' \otimes fst\ y' \mid_{rel$
 $snd\ x' \otimes snd\ y')$
proof–
have $(s' \otimes r \oplus s \otimes r', s \otimes s') \in carrier\ rel$
using *assms rel-def submonoid.m-closed*
by (*smt add.m-closed m-closed mem-Sigma-iff monoid.m-closed monoid-axioms*
partial-object.select-conv(1)
rev-subsetD subset)
have $(snd\ y' \otimes fst\ x' \oplus snd\ x' \otimes fst\ y', snd\ x' \otimes snd\ y') \in carrier\ rel$

```

using rel-def f1 f2 subset submonoid.m-closed eq-class-of-rng-of-frac-def
by (smt Product-Type.Collect-case-prodD add.m-closed mem-Sigma-iff mem-
ber-class-to-carrier
    partial-object.select-convs(1) semiring-simprules(3) rev-subsetD)
thus ?thesis
using elem-eq-class[of rel] equiv-obj-rng-of-frac
by (metis ⟨(s' ⊗ r ⊕ s ⊗ r', s ⊗ s') .=_rel (snd y' ⊗ fst x' ⊕ snd x' ⊗ fst y',
snd x' ⊗ snd y')⟩
    ⟨(s' ⊗ r ⊕ s ⊗ r', s ⊗ s') ∈ carrier rel⟩ class-of-to-rel)
qed
thus ?thesis
using f3
by simp
qed

```

lemma *closed-add-rng-of-frac*:

```

assumes (r, s) ∈ carrier rel and (r', s') ∈ carrier rel
shows (r |rel s) ⊕rec-rng-of-frac (r' |rel s') ∈ set-class-ofrel
proof –
have f1:(r |rel s) ⊕rec-rng-of-frac (r' |rel s') = (s' ⊗ r ⊕ s ⊗ r' |rel s ⊗ s')
using assms add-rng-of-frac-fundamental-lemma
by simp
have f2:s' ⊗ r ⊕ s ⊗ r' ∈ carrier R
using assms rel-def
by (metis (no-types, lifting) add.m-closed mem-Sigma-iff monoid.m-closed
monoid-axioms
    partial-object.select-convs(1) rev-subsetD subset)
have f3:s ⊗ s' ∈ S
using assms rel-def submonoid.m-closed
by simp
from f2 and f3 have (s' ⊗ r ⊕ s ⊗ r', s ⊗ s') ∈ carrier rel
by (simp add: rel-def)
thus ?thesis
using set-eq-class-of-rng-of-frac-def f1
by auto
qed

```

lemma *closed-rel-add*:

```

assumes (r, s) ∈ carrier rel and (r', s') ∈ carrier rel
shows (s' ⊗ r ⊕ s ⊗ r', s ⊗ s') ∈ carrier rel
proof –
have s ⊗ s' ∈ S
using assms rel-def submonoid.m-closed
by simp
have s' ⊗ r ⊕ s ⊗ r' ∈ carrier R
using assms rel-def
by (metis (no-types, lifting) add.m-closed mem-Sigma-iff monoid.m-closed
monoid-axioms
    partial-object.select-convs(1) rev-subsetD subset)

```

thus *?thesis*
using *rel-def*
by (*simp add: ‹s ⊗ s' ∈ S›*)
qed

lemma *assoc-add-rng-of-frac*:

assumes $(r, s) \in \text{carrier rel}$ **and** $(r', s') \in \text{carrier rel}$ **and** $(r'', s'') \in \text{carrier rel}$
shows $(r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r' \mid_{\text{rel}} s') \oplus_{\text{rec-rng-of-frac}} (r'' \mid_{\text{rel}} s'') =$
 $(r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} ((r' \mid_{\text{rel}} s') \oplus_{\text{rec-rng-of-frac}} (r'' \mid_{\text{rel}} s''))$

proof –

have $(r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r' \mid_{\text{rel}} s') = (s' \otimes r \oplus s \otimes r' \mid_{\text{rel}} s \otimes s')$
using *assms(1) assms(2) add-rng-of-frac-fundamental-lemma*
by *simp*
then have $f1: (r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r' \mid_{\text{rel}} s') \oplus_{\text{rec-rng-of-frac}} (r'' \mid_{\text{rel}} s'') =$
 $(s'' \otimes (s' \otimes r \oplus s \otimes r') \oplus (s \otimes s') \otimes r'' \mid_{\text{rel}} (s \otimes s') \otimes s'')$
using *assms add-rng-of-frac-fundamental-lemma closed-rel-add*
by *simp*
have $(r' \mid_{\text{rel}} s') \oplus_{\text{rec-rng-of-frac}} (r'' \mid_{\text{rel}} s'') = (s'' \otimes r' \oplus s' \otimes r'' \mid_{\text{rel}} s' \otimes s'')$
using *assms(2) assms(3) add-rng-of-frac-fundamental-lemma*
by *simp*
then have $f2: (r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} ((r' \mid_{\text{rel}} s') \oplus_{\text{rec-rng-of-frac}} (r'' \mid_{\text{rel}} s''))$
 $=$
 $((s' \otimes s'') \otimes r \oplus s \otimes (s'' \otimes r' \oplus s' \otimes r'') \mid_{\text{rel}} s \otimes (s' \otimes s''))$
using *assms add-rng-of-frac-fundamental-lemma closed-rel-add*
by *simp*
have $f3: (s \otimes s') \otimes s'' = s \otimes (s' \otimes s'')$
using *m-assoc subset assms rel-def*
by (*metis (no-types, lifting) mem-Sigma-iff partial-object.select-convs(1) rev-subsetD*)
have $s'' \otimes (s' \otimes r \oplus s \otimes r') \oplus (s \otimes s') \otimes r'' = (s' \otimes s'') \otimes r \oplus s \otimes (s'' \otimes r'$
 $\oplus s' \otimes r'')$
by (*smt a-assoc assms m-comm mem-Sigma-iff monoid.m-assoc monoid.m-closed monoid-axioms*
partial-object.select-convs(1) r-distr rel-def subset subset-iff)
thus *?thesis*
using *f1 f2 f3*
by *simp*
qed

lemma *add-rng-of-frac-zero*:

shows $(\mathbf{0} \mid_{\text{rel}} \mathbf{1}) \in \text{set-class-of rel}$
by (*metis (no-types, lifting) closed-mult-rng-of-frac mem-Sigma-iff monoid.simps(2) one-closed*
partial-object.select-convs(1) rec-monoid-rng-of-frac-def rel-def right-unit-mult-rng-of-frac semiring-simprules(4) zero-closed)

lemma *l-unit-add-rng-of-frac*:

assumes $(r, s) \in \text{carrier rel}$
shows $\mathbf{0}_{\text{rec-rng-of-frac}} \oplus_{\text{rec-rng-of-frac}} (r \mid_{\text{rel}} s) = (r \mid_{\text{rel}} s)$
proof –

have $(\mathbf{0} \mid_{rel} \mathbf{1}) \oplus_{rec-rng-of-frac} (r \mid_{rel} s) = (s \otimes \mathbf{0} \oplus \mathbf{1} \otimes r \mid_{rel} \mathbf{1} \otimes s)$
using *assms add-rng-of-frac-fundamental-lemma*
by *(simp add: rel-def)*
then have $(\mathbf{0} \mid_{rel} \mathbf{1}) \oplus_{rec-rng-of-frac} (r \mid_{rel} s) = (r \mid_{rel} s)$
using *assms rel-def subset*
by *auto*
thus *?thesis*
using *rec-rng-of-frac-def*
by *simp*
qed

lemma *r-unit-add-rng-of-frac:*
assumes $(r, s) \in carrier\ rel$
shows $(r \mid_{rel} s) \oplus_{rec-rng-of-frac} \mathbf{0}_{rec-rng-of-frac} = (r \mid_{rel} s)$
proof –
have $(r \mid_{rel} s) \oplus_{rec-rng-of-frac} (\mathbf{0} \mid_{rel} \mathbf{1}) = (\mathbf{1} \otimes r \oplus s \otimes \mathbf{0} \mid_{rel} s \otimes \mathbf{1})$
using *assms add-rng-of-frac-fundamental-lemma*
by *(simp add: rel-def)*
then have $(r \mid_{rel} s) \oplus_{rec-rng-of-frac} (\mathbf{0} \mid_{rel} \mathbf{1}) = (r \mid_{rel} s)$
using *assms rel-def subset*
by *auto*
thus *?thesis*
using *rec-rng-of-frac-def*
by *simp*
qed

lemma *comm-add-rng-of-frac:*
assumes $(r, s) \in carrier\ rel$ **and** $(r', s') \in carrier\ rel$
shows $(r \mid_{rel} s) \oplus_{rec-rng-of-frac} (r' \mid_{rel} s') = (r' \mid_{rel} s') \oplus_{rec-rng-of-frac} (r \mid_{rel} s)$
proof –
have $f1: (r \mid_{rel} s) \oplus_{rec-rng-of-frac} (r' \mid_{rel} s') = (s' \otimes r \oplus s \otimes r' \mid_{rel} s \otimes s')$
using *assms add-rng-of-frac-fundamental-lemma*
by *simp*
have $f2: (r' \mid_{rel} s') \oplus_{rec-rng-of-frac} (r \mid_{rel} s) = (s \otimes r' \oplus s' \otimes r \mid_{rel} s' \otimes s)$
using *assms add-rng-of-frac-fundamental-lemma*
by *simp*
thus *?thesis*
using *f1 f2*
by *(metis (no-types, lifting) add.m-comm assms(1) assms(2) m-comm mem-Sigma-iff*

partial-object.select-convs(1) rel-def semiring-simprules(3) rev-subsetD sub-
set)
qed

lemma *class-of-zero-rng-of-frac:*
assumes $s \in S$
shows $(\mathbf{0} \mid_{rel} s) = \mathbf{0}_{rec-rng-of-frac}$
proof –
have $f1: (\mathbf{0}, s) \in carrier\ rel$

using *assms rel-def*
by *simp*
have $\mathbf{1} \otimes (\mathbf{1} \otimes \mathbf{0} \ominus s \otimes \mathbf{0}) = \mathbf{0}$
using *assms local.ring-axioms rev-subsetD ring.ring-simprules(14) subset*
by *fastforce*
then have $(\mathbf{0}, s) \text{.}=\text{rel} (\mathbf{0}, \mathbf{1})$
using *rel-def submonoid.one-closed*
by *auto*
thus *?thesis*
using *elem-eq-class equiv-obj-rng-of-frac f1 rec-rng-of-frac-def*
by *(metis (no-types, lifting) class-of-to-rel mem-Sigma-iff one-closed partial-object.select-convs(1)*

rel-def ring-record-simps(11))
qed

lemma *r-inv-add-rng-of-frac*:
assumes $(r, s) \in \text{carrier } \text{rel}$
shows $(r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (\ominus r \mid_{\text{rel}} s) = \mathbf{0}_{\text{rec-rng-of-frac}}$
proof –
have $(\ominus r, s) \in \text{carrier } \text{rel}$
using *assms rel-def*
by *simp*
then have $(r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (\ominus r \mid_{\text{rel}} s) = (s \otimes r \oplus s \otimes \ominus r \mid_{\text{rel}} s \otimes s)$
using *assms add-rng-of-frac-fundamental-lemma*
by *simp*
then have $(r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (\ominus r \mid_{\text{rel}} s) = (\mathbf{0} \mid_{\text{rel}} s \otimes s)$
using *r-minus[of s r] assms rel-def subset rev-subsetD r-neg*
by *auto*
thus *?thesis*
using *class-of-zero-rng-of-frac assms rel-def submonoid.m-closed*
by *simp*
qed

lemma *l-inv-add-rng-of-frac*:
assumes $(r, s) \in \text{carrier } \text{rel}$
shows $(\ominus r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r \mid_{\text{rel}} s) = \mathbf{0}_{\text{rec-rng-of-frac}}$
proof –
have $(\ominus r, s) \in \text{carrier } \text{rel}$
using *assms rel-def*
by *simp*
then have $(\ominus r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r \mid_{\text{rel}} s) = (s \otimes \ominus r \oplus s \otimes r \mid_{\text{rel}} s \otimes s)$
using *assms add-rng-of-frac-fundamental-lemma*
by *simp*
then have $(\ominus r \mid_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r \mid_{\text{rel}} s) = (\mathbf{0} \mid_{\text{rel}} s \otimes s)$
using *r-minus[of s r] assms rel-def subset rev-subsetD l-neg*
by *auto*
thus *?thesis*
using *class-of-zero-rng-of-frac assms rel-def submonoid.m-closed*
by *simp*

qed

lemma *abelian-group-rng-of-frac*:

shows *abelian-group (rec-rng-of-frac)*

proof

show $\bigwedge x y. \llbracket x \in \text{carrier (add-monoid rec-rng-of-frac)};$
 $y \in \text{carrier (add-monoid rec-rng-of-frac)} \rrbracket$

$\implies x \otimes_{\text{add-monoid rec-rng-of-frac}} y$
 $\in \text{carrier (add-monoid rec-rng-of-frac)}$

using *closed-add-rng-of-frac*

by (*smt mem-Collect-eq monoid.select-convs(1) partial-object.select-convs(1)*)

rec-rng-of-frac-def

set-eq-class-of-rng-of-frac-def)

show $\bigwedge x y z.$

$\llbracket x \in \text{carrier (add-monoid rec-rng-of-frac)};$
 $y \in \text{carrier (add-monoid rec-rng-of-frac)};$
 $z \in \text{carrier (add-monoid rec-rng-of-frac)} \rrbracket$

$\implies x \otimes_{\text{add-monoid rec-rng-of-frac}} y \otimes_{\text{add-monoid rec-rng-of-frac}} z =$
 $x \otimes_{\text{add-monoid rec-rng-of-frac}} (y \otimes_{\text{add-monoid rec-rng-of-frac}} z)$

using *assoc-add-rng-of-frac*

by (*smt mem-Collect-eq monoid.simps(1) partial-object.select-convs(1) rec-rng-of-frac-def*

set-eq-class-of-rng-of-frac-def)

show $\mathbf{1}_{\text{add-monoid rec-rng-of-frac}} \in \text{carrier (add-monoid rec-rng-of-frac)}$

using *add-rng-of-frac-zero* **by** (*simp add: rec-rng-of-frac-def*)

show $\bigwedge x. x \in \text{carrier (add-monoid rec-rng-of-frac)} \implies$

$\mathbf{1}_{\text{add-monoid rec-rng-of-frac}} \otimes_{\text{add-monoid rec-rng-of-frac}} x = x$

using *l-unit-add-rng-of-frac*

by (*smt mem-Collect-eq monoid.select-convs(1) monoid.select-convs(2) partial-object.select-convs(1)*

rec-rng-of-frac-def set-eq-class-of-rng-of-frac-def)

show $\bigwedge x. x \in \text{carrier (add-monoid rec-rng-of-frac)} \implies$

$x \otimes_{\text{add-monoid rec-rng-of-frac}} \mathbf{1}_{\text{add-monoid rec-rng-of-frac}} = x$

using *r-unit-add-rng-of-frac*

by (*smt mem-Collect-eq monoid.select-convs(1) monoid.select-convs(2) partial-object.select-convs(1)*

rec-rng-of-frac-def set-eq-class-of-rng-of-frac-def)

show $\bigwedge x y. \llbracket x \in \text{carrier (add-monoid rec-rng-of-frac)};$ $y \in \text{carrier (add-monoid rec-rng-of-frac)} \rrbracket$

$\implies x \otimes_{\text{add-monoid rec-rng-of-frac}} y = y \otimes_{\text{add-monoid rec-rng-of-frac}} x$

using *comm-add-rng-of-frac*

by (*smt mem-Collect-eq monoid.select-convs(1) partial-object.select-convs(1)*)

rec-rng-of-frac-def

set-eq-class-of-rng-of-frac-def)

show $\text{carrier (add-monoid rec-rng-of-frac)} \subseteq \text{Units (add-monoid rec-rng-of-frac)}$

proof

show $x \in \text{Units (add-monoid rec-rng-of-frac)}$ **if** $x \in \text{carrier (add-monoid rec-rng-of-frac)}$ **for** x

proof –

have $x \in \text{set-class-of}_{rel}$
using *that rec-rng-of-frac-def by simp*
then obtain r **and** s **where** $f1:(r, s) \in \text{carrier } rel$ **and** $f2:x = (r \mid_{rel} s)$
using *set-eq-class-of-rng-of-frac-def*
by *(smt mem-Collect-eq)*
then have $f3:(r \mid_{rel} s) \oplus_{rec-rng-of-frac} (\ominus r \mid_{rel} s) = \mathbf{0}_{rec-rng-of-frac}$
using $f1$ *r-inv-add-rng-of-frac[of r s]*
by *simp*
have $f4:(\ominus r \mid_{rel} s) \oplus_{rec-rng-of-frac} (r \mid_{rel} s) = \mathbf{0}_{rec-rng-of-frac}$
using $f1$ *l-inv-add-rng-of-frac[of r s]*
by *simp*
then have $\exists y \in \text{set-class-of}_{rel}. y \oplus_{rec-rng-of-frac} x = \mathbf{0}_{rec-rng-of-frac} \wedge x$
 $\oplus_{rec-rng-of-frac} y = \mathbf{0}_{rec-rng-of-frac}$
using $f2$ $f3$ $f4$
by *(metis (no-types, lifting) abelian-group.a-inv-closed class-of-zero-rng-of-frac*

closed-add-rng-of-frac f1 is-abelian-group mem-Sigma-iff partial-object.select-convs(1)

rel-def r-unit-add-rng-of-frac zero-closed)
thus $x \in \text{Units}$ *(add-monoid rec-rng-of-frac)*
using *rec-rng-of-frac-def that by (simp add: Units-def)*
qed
qed
qed

lemma *r-distr-rng-of-frac:*
assumes $(r, s) \in \text{carrier } rel$ **and** $(r', s') \in \text{carrier } rel$ **and** $(r'', s'') \in \text{carrier } rel$
shows $((r \mid_{rel} s) \oplus_{rec-rng-of-frac} (r' \mid_{rel} s')) \otimes_{rec-rng-of-frac} (r'' \mid_{rel} s'') =$
 $(r \mid_{rel} s) \otimes_{rec-rng-of-frac} (r'' \mid_{rel} s'') \oplus_{rec-rng-of-frac} (r' \mid_{rel} s') \otimes_{rec-rng-of-frac}$
 $(r'' \mid_{rel} s'')$
proof –
have $(r \mid_{rel} s) \oplus_{rec-rng-of-frac} (r' \mid_{rel} s') = (s' \otimes r \oplus s \otimes r' \mid_{rel} s \otimes s')$
using $assms(1)$ $assms(2)$ *add-rng-of-frac-fundamental-lemma*
by *simp*
then have $f1:((r \mid_{rel} s) \oplus_{rec-rng-of-frac} (r' \mid_{rel} s')) \otimes_{rec-rng-of-frac} (r'' \mid_{rel} s'')$
 $=$
 $((s' \otimes r \oplus s \otimes r') \otimes r'' \mid_{rel} (s \otimes s') \otimes s'')$
using $assms$ *mult-rng-of-frac-fundamental-lemma*
by *(simp add: closed-rel-add rec-monoid-rng-of-frac-def rec-rng-of-frac-def)*
have $f2:(r \mid_{rel} s) \otimes_{rec-rng-of-frac} (r'' \mid_{rel} s'') = (r \otimes r'' \mid_{rel} s \otimes s'')$
using $assms(1)$ $assms(3)$ *mult-rng-of-frac-fundamental-lemma*
by *(simp add: rec-monoid-rng-of-frac-def rec-rng-of-frac-def)*
have $f3:(r' \mid_{rel} s') \otimes_{rec-rng-of-frac} (r'' \mid_{rel} s'') = (r' \otimes r'' \mid_{rel} s' \otimes s'')$
using $assms(2)$ $assms(3)$ *mult-rng-of-frac-fundamental-lemma*
by *(simp add: rec-monoid-rng-of-frac-def rec-rng-of-frac-def)*
have $f4:(r \otimes r'', s \otimes s'') \in \text{carrier } rel$
using *rel-def assms(1) assms(3) submonoid.m-closed*
by *simp*
have $f5:(r' \otimes r'', s' \otimes s'') \in \text{carrier } rel$

using *rel-def* *assms(2)* *assms(3)* *submonoid.m-closed*
by *simp*
from *f2* **and** *f3* **have** *f6*: $(r \mid_{\text{rel}} s) \otimes_{\text{rec-rng-of-frac}} (r'' \mid_{\text{rel}} s'') \oplus_{\text{rec-rng-of-frac}} (r' \mid_{\text{rel}} s') \otimes_{\text{rec-rng-of-frac}} (r'' \mid_{\text{rel}} s'')$
 $= ((s' \otimes s'') \otimes (r \otimes r'') \oplus (s \otimes s') \otimes (r' \otimes r'')) \mid_{\text{rel}} (s \otimes s'') \otimes (s' \otimes s')$
using *assms f4 f5 submonoid.m-closed add-rng-of-frac-fundamental-lemma*
by *simp*
have $(s \otimes s'' \otimes (s' \otimes s')) \otimes ((s' \otimes r \oplus s \otimes r') \otimes r'') = (s \otimes s'' \otimes (s' \otimes s'))$
 $\otimes (s' \otimes r \otimes r'' \oplus s \otimes r' \otimes r'')$
using *assms rel-def subset rev-subsetD l-distr*
by (*smt mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1)*)
then have *f7*: $(s \otimes s'' \otimes (s' \otimes s')) \otimes ((s' \otimes r \oplus s \otimes r') \otimes r'') =$
 $(s \otimes s'' \otimes (s' \otimes s')) \otimes (s' \otimes r \otimes r'') \oplus (s \otimes s'' \otimes (s' \otimes s')) \otimes (s \otimes r' \otimes r'')$
using *assms rel-def subset rev-subsetD submonoid.m-closed r-distr*
by (*smt mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1)*)
have *f8*: $(s \otimes s' \otimes s'') \otimes (s' \otimes s'' \otimes (r \otimes r'')) \oplus s \otimes s'' \otimes (r' \otimes r'') =$
 $(s \otimes s' \otimes s'') \otimes (s' \otimes s'' \otimes (r \otimes r'')) \oplus (s \otimes s' \otimes s'') \otimes (s \otimes s'' \otimes (r' \otimes r''))$
using *assms rel-def subset rev-subsetD submonoid.m-closed r-distr*
by (*smt mem-Sigma-iff partial-object.select-convs(1) semiring-simprules(3)*)
have $(s \otimes s'' \otimes (s' \otimes s')) = (s \otimes (s'' \otimes s') \otimes s')$
using *assms rel-def subset rev-subsetD submonoid.m-closed m-assoc*
by (*smt mem-Sigma-iff partial-object.select-convs(1) semiring-simprules(3)*)
then have *f9*: $(s \otimes s'' \otimes (s' \otimes s')) = (s \otimes s' \otimes (s'' \otimes s'))$
using *assms rel-def subset rev-subsetD submonoid.m-closed m-comm m-assoc*
by (*smt mem-Sigma-iff partial-object.select-convs(1) semiring-simprules(3)*)
then have *f10*: $(s \otimes s'' \otimes (s' \otimes s')) \otimes (s' \otimes r \otimes r'') = (s \otimes s' \otimes s'') \otimes (s' \otimes$
 $s'' \otimes (r \otimes r''))$
using *assms rel-def subset rev-subsetD submonoid.m-closed m-assoc m-comm*
by (*smt mem-Sigma-iff partial-object.select-convs(1) semiring-simprules(3)*)
have $(s \otimes s'' \otimes (r' \otimes r'')) = (s'' \otimes s \otimes (r' \otimes r''))$
using *assms rel-def subset rev-subsetD m-comm*
by (*metis (no-types, lifting) mem-Sigma-iff partial-object.select-convs(1)*)
then have $(s \otimes s'' \otimes (s' \otimes s')) \otimes (s \otimes r' \otimes r'') = (s \otimes s' \otimes s'') \otimes (s \otimes s'' \otimes$
 $(r' \otimes r''))$
using *assms rel-def subset rev-subsetD submonoid.m-closed m-comm m-assoc f9*
by (*smt mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1)*)
then have $((s \otimes s'' \otimes (s' \otimes s')) \otimes ((s' \otimes r \oplus s \otimes r') \otimes r'')) = (s \otimes s' \otimes s'')$
 $\otimes (s' \otimes s'' \otimes (r \otimes r'')) \oplus s \otimes s'' \otimes (r' \otimes r''))$
using *f7 f8 f10*
by *presburger*
then have $((s \otimes s'' \otimes (s' \otimes s')) \otimes ((s' \otimes r \oplus s \otimes r') \otimes r'')) \ominus (s \otimes s' \otimes s'')$
 $\otimes (s' \otimes s'' \otimes (r \otimes r'')) \oplus s \otimes s'' \otimes (r' \otimes r'')) = \mathbf{0}$
by (*smt a-minus-def assms(1) assms(2) assms(3) closed-rel-add mem-Sigma-iff*
partial-object.select-convs(1) r-neg rel-def semiring-simprules(3) rev-subsetD
subset)
then have *f11*: $\mathbf{1} \otimes (((s \otimes s'' \otimes (s' \otimes s')) \otimes ((s' \otimes r \oplus s \otimes r') \otimes r'')) \ominus (s \otimes$
 $s' \otimes s'') \otimes (s' \otimes s'' \otimes (r \otimes r'')) \oplus s \otimes s'' \otimes (r' \otimes r'')))) = \mathbf{0}$
by *simp*

have $f12:((s' \otimes r \oplus s \otimes r') \otimes r'', s \otimes s' \otimes s'') \in \text{carrier rel}$
using *assms closed-rel-add rel-def*
by *auto*
have $f13:(s' \otimes s'' \otimes (r \otimes r'') \oplus s \otimes s'' \otimes (r' \otimes r''), s \otimes s'' \otimes (s' \otimes s'')) \in$
carrier rel
by (*simp add: closed-rel-add f4 f5*)
have $1 \in S$
using *submonoid.one-closed*
by *simp*
then have $((s' \otimes r \oplus s \otimes r') \otimes r'', s \otimes s' \otimes s'') \dot{=}_{\text{rel}} (s' \otimes s'' \otimes (r \otimes r'') \oplus$
 $s \otimes s'' \otimes (r' \otimes r''), s \otimes s'' \otimes (s' \otimes s''))$
using *rel-def f11 f13 f12*
by *auto*
then have $((s' \otimes r \oplus s \otimes r') \otimes r'' |_{\text{rel}} s \otimes s' \otimes s'') = (s' \otimes s'' \otimes (r \otimes r'') \oplus s$
 $\otimes s'' \otimes (r' \otimes r'') |_{\text{rel}} s \otimes s'' \otimes (s' \otimes s''))$
using *elem-eq-class*
by (*metis class-of-to-rel equiv-obj-rng-of-frac f12 f13*)
thus *?thesis*
using *f1 f6*
by *simp*
qed

lemma *l-distr-rng-of-frac:*

assumes $(r, s) \in \text{carrier rel}$ **and** $(r', s') \in \text{carrier rel}$ **and** $(r'', s'') \in \text{carrier rel}$
shows $(r'' |_{\text{rel}} s'') \otimes_{\text{rec-rng-of-frac}} ((r |_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r' |_{\text{rel}} s')) =$
 $(r'' |_{\text{rel}} s'') \otimes_{\text{rec-rng-of-frac}} (r |_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r'' |_{\text{rel}} s'') \otimes_{\text{rec-rng-of-frac}}$
 $(r' |_{\text{rel}} s')$

proof –

have $(r |_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r' |_{\text{rel}} s') = (s' \otimes r \oplus s \otimes r' |_{\text{rel}} s \otimes s')$
using *assms(1) assms(2) add-rng-of-frac-fundamental-lemma*
by *simp*
then have $f1:(r'' |_{\text{rel}} s'') \otimes_{\text{rec-rng-of-frac}} ((r |_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}} (r' |_{\text{rel}} s'))$
 $=$
 $(r'' \otimes (s' \otimes r \oplus s \otimes r') |_{\text{rel}} s'' \otimes (s \otimes s'))$
using *assms mult-rng-of-frac-fundamental-lemma*
by (*simp add: closed-rel-add rec-monoid-rng-of-frac-def rec-rng-of-frac-def*)
have $f2:(r'' |_{\text{rel}} s'') \otimes_{\text{rec-rng-of-frac}} (r |_{\text{rel}} s) = (r'' \otimes r |_{\text{rel}} s'' \otimes s)$
using *assms(1) assms(3) mult-rng-of-frac-fundamental-lemma*
by (*simp add: rec-monoid-rng-of-frac-def rec-rng-of-frac-def*)
have $f3:(r'' |_{\text{rel}} s'') \otimes_{\text{rec-rng-of-frac}} (r' |_{\text{rel}} s') = (r'' \otimes r' |_{\text{rel}} s'' \otimes s')$
using *assms(2) assms(3) mult-rng-of-frac-fundamental-lemma*
by (*simp add: rec-monoid-rng-of-frac-def rec-rng-of-frac-def*)
have $f4:(r'' \otimes r, s'' \otimes s) \in \text{carrier rel}$
using *rel-def assms(1) assms(3) submonoid.m-closed*
by *simp*
have $f5:(r'' \otimes r', s'' \otimes s') \in \text{carrier rel}$
using *rel-def assms(2) assms(3) submonoid.m-closed*
by *simp*
from $f2$ **and** $f3$ **have** $f6:(r'' |_{\text{rel}} s'') \otimes_{\text{rec-rng-of-frac}} (r |_{\text{rel}} s) \oplus_{\text{rec-rng-of-frac}}$

$(r'' \mid_{rel} s'') \otimes_{rec-rng-of-frac} (r' \mid_{rel} s')$
 $= ((s'' \otimes s') \otimes (r'' \otimes r) \oplus (s'' \otimes s) \otimes (r'' \otimes r')) \mid_{rel} (s'' \otimes s) \otimes (s'' \otimes s')$
using *assms f4 f5 submonoid.m-closed add-rng-of-frac-fundamental-lemma*
by *simp*
have $(s'' \otimes s \otimes (s'' \otimes s')) \otimes (r'' \otimes (s' \otimes r \oplus s \otimes r')) = (s'' \otimes s \otimes (s'' \otimes s'))$
 $\otimes (r'' \otimes (s' \otimes r) \oplus r'' \otimes (s \otimes r'))$
using *assms rel-def subset rev-subsetD r-distr*
by *(smt mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1))*
then have $f7:(s'' \otimes s \otimes (s'' \otimes s')) \otimes (r'' \otimes (s' \otimes r \oplus s \otimes r')) =$
 $(s'' \otimes s \otimes (s'' \otimes s')) \otimes (r'' \otimes (s' \otimes r)) \oplus (s'' \otimes s \otimes (s'' \otimes s')) \otimes (r'' \otimes (s \otimes$
 $r'))$
using *assms rel-def subset rev-subsetD submonoid.m-closed r-distr*
by *(smt mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1))*
have $f8:(s'' \otimes s \otimes s') \otimes (s'' \otimes s' \otimes (r'' \otimes r) \oplus s'' \otimes s \otimes (r'' \otimes r')) =$
 $(s'' \otimes s \otimes s') \otimes (s'' \otimes s' \otimes (r'' \otimes r)) \oplus (s'' \otimes s \otimes s') \otimes (s'' \otimes s \otimes (r'' \otimes r'))$
using *assms rel-def subset rev-subsetD submonoid.m-closed r-distr*
by *(smt mem-Sigma-iff partial-object.select-convs(1) semiring-simprules(3))*
have $(s'' \otimes s \otimes (s'' \otimes s')) = (s'' \otimes (s \otimes s')) \otimes s'$
using *assms rel-def subset rev-subsetD submonoid.m-closed m-assoc*
by *(smt mem-Sigma-iff partial-object.select-convs(1) semiring-simprules(3))*
then have $f9:(s'' \otimes s \otimes (s'' \otimes s')) = (s'' \otimes s'' \otimes (s \otimes s'))$
using *assms rel-def subset rev-subsetD submonoid.m-closed m-comm m-assoc*
by *(smt mem-Sigma-iff partial-object.select-convs(1) semiring-simprules(3))*
then have $f10:(s'' \otimes s \otimes (s'' \otimes s')) \otimes (r'' \otimes s' \otimes r) = (s'' \otimes s \otimes s') \otimes (s'' \otimes$
 $s' \otimes (r'' \otimes r))$
using *assms rel-def subset rev-subsetD submonoid.m-closed m-assoc m-comm*
by *(smt mem-Sigma-iff partial-object.select-convs(1) semiring-simprules(3))*
have $(s'' \otimes s \otimes (r'' \otimes r')) = (s \otimes s'' \otimes (r'' \otimes r'))$
using *assms rel-def subset rev-subsetD m-comm*
by *(metis (no-types, lifting) mem-Sigma-iff partial-object.select-convs(1))*
then have $(s'' \otimes s \otimes (s'' \otimes s')) \otimes (r'' \otimes s \otimes r') = (s'' \otimes s \otimes s') \otimes (s'' \otimes s \otimes$
 $(r'' \otimes r'))$
using *assms rel-def subset rev-subsetD submonoid.m-closed m-comm m-assoc f9*
by *(smt mem-Sigma-iff monoid.m-closed monoid-axioms partial-object.select-convs(1))*
then have $((s'' \otimes s \otimes (s'' \otimes s')) \otimes (r'' \otimes (s' \otimes r \oplus s \otimes r'))) = (s'' \otimes (s \otimes s'))$
 $\otimes (s'' \otimes s' \otimes (r'' \otimes r) \oplus s'' \otimes s \otimes (r'' \otimes r'))$
using *f7 f8 f10*
by *(smt assms(1) assms(2) assms(3) m-assoc mem-Sigma-iff partial-object.select-convs(1) rel-def*
rev-subsetD subset)
then have $((s'' \otimes s \otimes (s'' \otimes s')) \otimes (r'' \otimes (s' \otimes r \oplus s \otimes r'))) \ominus (s'' \otimes (s \otimes s'))$
 $\otimes (s'' \otimes s' \otimes (r'' \otimes r) \oplus s'' \otimes s \otimes (r'' \otimes r')) = \mathbf{0}$
by *(smt a-minus-def assms(1) assms(2) assms(3) closed-rel-add mem-Sigma-iff*
partial-object.select-convs(1)
r-neg rel-def semiring-simprules(3) rev-subsetD subset)
then have $f11:\mathbf{1} \otimes (((s'' \otimes s \otimes (s'' \otimes s')) \otimes (r'' \otimes (s' \otimes r \oplus s \otimes r'))) \ominus (s'' \otimes$
 $(s \otimes s')) \otimes (s'' \otimes s' \otimes (r'' \otimes r) \oplus s'' \otimes s \otimes (r'' \otimes r')))) = \mathbf{0}$
by *simp*
have $f12:(r'' \otimes (s' \otimes r \oplus s \otimes r'), s'' \otimes (s \otimes s')) \in carrier\ rel$

using *assms closed-rel-add rel-def*
by *auto*
have $f13:(s'' \otimes s' \otimes (r'' \otimes r) \oplus s'' \otimes s \otimes (r'' \otimes r'), s'' \otimes s \otimes (s'' \otimes s')) \in$
carrier rel
by (*simp add: closed-rel-add f4 f5*)
have $1 \in S$
using *submonoid.one-closed*
by *simp*
then have $(r'' \otimes (s' \otimes r \oplus s \otimes r'), s'' \otimes (s \otimes s')) \dot{=}_{rel} (s'' \otimes s' \otimes (r'' \otimes r)$
 $\oplus s'' \otimes s \otimes (r'' \otimes r'), s'' \otimes s \otimes (s'' \otimes s'))$
using *rel-def f11 f13 f12*
by *auto*
then have $(r'' \otimes (s' \otimes r \oplus s \otimes r') \mid_{rel} s'' \otimes (s \otimes s')) = (s'' \otimes s' \otimes (r'' \otimes r)$
 $\oplus s'' \otimes s \otimes (r'' \otimes r') \mid_{rel} s'' \otimes s \otimes (s'' \otimes s'))$
using *elem-eq-class*
by (*metis class-of-to-rel equiv-obj-rng-of-frac f12 f13*)
thus *?thesis*
using *f1 f6*
by *simp*
qed

lemma *rng-rng-of-frac*:
shows *ring (rec-rng-of-frac)*
proof –
have $f1:\forall x y z. x \in \text{carrier } \text{rec-rng-of-frac} \longrightarrow y \in \text{carrier } \text{rec-rng-of-frac} \longrightarrow z$
 $\in \text{carrier } \text{rec-rng-of-frac}$
 $\longrightarrow (x \oplus_{\text{rec-rng-of-frac}} y) \otimes_{\text{rec-rng-of-frac}} z = x \otimes_{\text{rec-rng-of-frac}} z \oplus_{\text{rec-rng-of-frac}}$
 $y \otimes_{\text{rec-rng-of-frac}} z$
using *r-distr-rng-of-frac rec-rng-of-frac-def*
by (*smt mem-Collect-eq partial-object.select-convs(1) set-eq-class-of-rng-of-frac-def*)
have $f2:\forall x y z. x \in \text{carrier } \text{rec-rng-of-frac} \longrightarrow y \in \text{carrier } \text{rec-rng-of-frac} \longrightarrow z$
 $\in \text{carrier } \text{rec-rng-of-frac}$
 $\longrightarrow z \otimes_{\text{rec-rng-of-frac}} (x \oplus_{\text{rec-rng-of-frac}} y) = z \otimes_{\text{rec-rng-of-frac}} x \oplus_{\text{rec-rng-of-frac}}$
 $z \otimes_{\text{rec-rng-of-frac}} y$
using *l-distr-rng-of-frac rec-rng-of-frac-def*
by (*smt mem-Collect-eq partial-object.select-convs(1) set-eq-class-of-rng-of-frac-def*)
then have *ring-axioms (rec-rng-of-frac)*
using *ring-axioms-def f1 f2*
by *auto*
thus *?thesis*
using *ring-def[of rec-rng-of-frac] abelian-group-rng-of-frac monoid-rng-of-frac*
rec-rng-of-frac-def
abelian-group-axioms-def rec-monoid-rng-of-frac-def eq-class-of-rng-of-frac-def
by (*simp add: Group.monoid-def*)
qed

lemma *crng-rng-of-frac*:
shows *cring (rec-rng-of-frac)*
using *cring-def[of rec-rng-of-frac] rng-rng-of-frac comm-monoid-rng-of-frac rec-rng-of-frac-def*

rec-monoid-rng-of-frac-def eq-class-of-rng-of-frac-def
by (*metis (no-types, lifting) comm-monoid.m-comm monoid.monoid-comm-monoidI*
monoid.select-convs(1)
partial-object.select-convs(1) ring.is-monoid)

lemma *simp-in-frac*:

assumes $(r, s) \in \text{carrier } \text{rel}$ **and** $s' \in S$
shows $(r \mid_{\text{rel}} s) = (s' \otimes r \mid_{\text{rel}} s' \otimes s)$

proof –

have $f1: (s' \otimes r, s' \otimes s) \in \text{carrier } \text{rel}$
using *assms rel-def submonoid.m-closed subset rev-subsetD*
by *auto*
have $(s' \otimes s) \otimes r \ominus s \otimes (s' \otimes r) = (s' \otimes s) \otimes r \ominus (s \otimes s') \otimes r$
using *assms subset rev-subsetD m-assoc[of s s' r] rel-def*
by (*metis (no-types, lifting) mem-Sigma-iff partial-object.select-convs(1)*)
then have $(s' \otimes s) \otimes r \ominus s \otimes (s' \otimes r) = (s' \otimes s) \otimes r \ominus (s' \otimes s) \otimes r$
using *m-comm[of s s'] assms subset rev-subsetD rel-def*
by (*metis (no-types, lifting) mem-Sigma-iff partial-object.select-convs(1)*)
then have $(s' \otimes s) \otimes r \ominus s \otimes (s' \otimes r) = \mathbf{0}$
by (*metis (no-types, lifting) a-minus-def assms mem-Sigma-iff partial-object.select-convs(1)*)

r-neg rel-def semiring-simprules(3) rev-subsetD subset
then have $\mathbf{1} \otimes ((s' \otimes s) \otimes r \ominus s \otimes (s' \otimes r)) = \mathbf{0}$
by *simp*
then have $(r, s) \text{.}=\text{rel} (s' \otimes r, s' \otimes s)$
using *assms(1) f1 rel-def one-closed*
by *auto*
thus *?thesis*
using *elem-eq-class*
by (*metis assms(1) class-of-to-rel equiv-obj-rng-of-frac f1*)

qed

1.2 The Natural Homomorphism from a Ring to Its Localization

definition *rng-to-rng-of-frac* :: $'a \Rightarrow ('a \times 'a)$ *set where*
rng-to-rng-of-frac $r \equiv (r \mid_{\text{rel}} \mathbf{1})$

lemma *rng-to-rng-of-frac-is-ring-hom* :

shows *rng-to-rng-of-frac* \in *ring-hom* R *rec-rng-of-frac*

proof –

have $f1: \text{rng-to-rng-of-frac} \in \text{carrier } R \rightarrow \text{carrier } \text{rec-rng-of-frac}$
using *rng-to-rng-of-frac-def rec-rng-of-frac-def set-eq-class-of-rng-of-frac-def*
rel-def
by *fastforce*
have $f2: \forall x y. x \in \text{carrier } R \wedge y \in \text{carrier } R$
 $\longrightarrow \text{rng-to-rng-of-frac} (x \otimes_R y) = \text{rng-to-rng-of-frac } x \otimes_{\text{rec-rng-of-frac}} \text{rng-to-rng-of-frac } y$

```

proof(rule allI, rule allI, rule impI)
  fix x y
  assume x ∈ carrier R ∧ y ∈ carrier R
  have f1: rng-to-rng-of-frac (x ⊗R y) = (x ⊗ y |rel 1)
    using rng-to-rng-of-frac-def
    by simp
  have rng-to-rng-of-frac x ⊗rec-rng-of-frac rng-to-rng-of-frac y = (x |rel 1)
    ⊗rec-rng-of-frac (y |rel 1)
    using rng-to-rng-of-frac-def
    by simp
  then have rng-to-rng-of-frac x ⊗rec-rng-of-frac rng-to-rng-of-frac y = (x ⊗ y
|rel 1)
    using mult-rng-of-frac-fundamental-lemma
    by (simp add: ⟨x ∈ carrier R ∧ y ∈ carrier R⟩ rec-monoid-rng-of-frac-def
rec-rng-of-frac-def rel-def)
  thus rng-to-rng-of-frac (x ⊗R y) = rng-to-rng-of-frac x ⊗rec-rng-of-frac rng-to-rng-of-frac
y
    using f1
    by auto
qed
have f3: ∀ x y. x ∈ carrier R ∧ y ∈ carrier R
  → rng-to-rng-of-frac (x ⊕R y) = rng-to-rng-of-frac x ⊕rec-rng-of-frac rng-to-rng-of-frac
y
proof(rule allI, rule allI, rule impI)
  fix x y
  assume a: x ∈ carrier R ∧ y ∈ carrier R
  have f1: rng-to-rng-of-frac (x ⊕R y) = (x ⊕ y |rel 1)
    using rng-to-rng-of-frac-def
    by simp
  have rng-to-rng-of-frac x ⊕rec-rng-of-frac rng-to-rng-of-frac y = (x |rel 1)
    ⊕rec-rng-of-frac (y |rel 1)
    using rng-to-rng-of-frac-def
    by simp
  then have rng-to-rng-of-frac x ⊕rec-rng-of-frac rng-to-rng-of-frac y = (1 ⊗ x
⊕ 1 ⊗ y |rel 1 ⊗ 1)
    using mult-rng-of-frac-fundamental-lemma a
    eq-obj-rng-of-frac.add-rng-of-frac-fundamental-lemma eq-obj-rng-of-frac.rng-to-rng-of-frac-def

    eq-obj-rng-of-frac-axioms f1
    by fastforce
  then have rng-to-rng-of-frac x ⊕rec-rng-of-frac rng-to-rng-of-frac y = (x ⊕ y
|rel 1)
    using l-one a
    by simp
  thus rng-to-rng-of-frac (x ⊕R y) = rng-to-rng-of-frac x ⊕rec-rng-of-frac rng-to-rng-of-frac
y
    using f1
    by auto
qed

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have rng-to-rng-of-frac 1 = (1 |rel 1)
  using rng-to-rng-of-frac-def
  by simp
then have rng-to-rng-of-frac 1R = 1rec-rng-of-frac
  using rec-rng-of-frac-def
  by simp
thus ?thesis
  using ring-hom-def[of R rec-rng-of-frac] f1 f2 f3 f4
  by simp
qed

lemma Im-rng-to-rng-of-frac-unit:
  assumes  $x \in \text{rng-to-rng-of-frac } S$ 
  shows  $x \in \text{Units rec-rng-of-frac}$ 
proof –
  obtain  $s$  where  $a1:s \in S$  and  $a2:x = (s \text{ |}_{rel} \mathbf{1})$ 
    using assms rng-to-rng-of-frac-def rel-def
    by auto
  then have  $(s \text{ |}_{rel} \mathbf{1}) \otimes_{rec-rng-of-frac} (\mathbf{1} \text{ |}_{rel} s) = (s \otimes \mathbf{1} \text{ |}_{rel} s \otimes \mathbf{1})$ 
    using mult-rng-of-frac-fundamental-lemma rec-monoid-rng-of-frac-def rec-rng-of-frac-def
  rel-def subset
    by auto
  then have  $f1:(s \text{ |}_{rel} \mathbf{1}) \otimes_{rec-rng-of-frac} (\mathbf{1} \text{ |}_{rel} s) = (\mathbf{1} \text{ |}_{rel} \mathbf{1})$ 
    using simp-in-frac a1 rel-def
    by auto
  have  $(\mathbf{1} \text{ |}_{rel} s) \otimes_{rec-rng-of-frac} (s \text{ |}_{rel} \mathbf{1}) = (s \otimes \mathbf{1} \text{ |}_{rel} s \otimes \mathbf{1})$ 
    using mult-rng-of-frac-fundamental-lemma rec-monoid-rng-of-frac-def rec-rng-of-frac-def
  rel-def
    subset a1
    by auto
  then have  $f2:(\mathbf{1} \text{ |}_{rel} s) \otimes_{rec-rng-of-frac} (s \text{ |}_{rel} \mathbf{1}) = (\mathbf{1} \text{ |}_{rel} \mathbf{1})$ 
    using simp-in-frac a1 rel-def
    by auto
  then have  $f3:\exists y \in \text{carrier rec-rng-of-frac. } y \otimes_{rec-rng-of-frac} x = \mathbf{1}_{rec-rng-of-frac}$ 
   $\wedge$ 
     $x \otimes_{rec-rng-of-frac} y = \mathbf{1}_{rec-rng-of-frac}$ 
    using rec-rng-of-frac-def f1 f2 a2 rel-def a1
  by (metis (no-types, lifting) class-of-zero-rng-of-frac closed-add-rng-of-frac l-unit-add-rng-of-frac
    mem-Sigma-iff monoid.select-convs(2) partial-object.select-convs(1) semir-
ing-simprules(4) zero-closed)
  have  $x \in \text{carrier rec-rng-of-frac}$ 
    using a2 a1 subset rev-subsetD rec-rng-of-frac-def
  by (metis (no-types, opaque-lifting) ring-hom-closed rng-to-rng-of-frac-def rng-to-rng-of-frac-is-ring-hom)
  thus ?thesis
    using Units-def[of rec-rng-of-frac] f3
    by auto
qed

```

lemma *eq-class-to-rel*:

assumes $(r, s) \in \text{carrier } R \times S$ **and** $(r', s') \in \text{carrier } R \times S$ **and** $(r \mid_{\text{rel}} s) = (r' \mid_{\text{rel}} s')$

shows $(r, s) \text{.}=\text{rel } (r', s')$

proof –

have $(r, s) \in (r \mid_{\text{rel}} s)$

using *assms(1) equiv-obj-rng-of-frac equivalence-def*

by (*metis (no-types, lifting) CollectI case-prodI eq-class-of-rng-of-frac-def partial-object.select-convs(1) rel-def*)

then have $(r, s) \in (r' \mid_{\text{rel}} s')$

using *assms(3)*

by *simp*

then have $(r', s') \text{.}=\text{rel } (r, s)$

by (*simp add: eq-class-of-rng-of-frac-def*)

thus *?thesis*

using *equiv-obj-rng-of-frac equivalence-def*

by (*metis (no-types, lifting) assms(1) assms(2) partial-object.select-convs(1) rel-def*)

qed

lemma *rng-to-rng-of-frac-without-zero-div-is-inj*:

assumes $0 \notin S$ **and** $\forall a \in \text{carrier } R. \forall b \in \text{carrier } R. a \otimes b = 0 \longrightarrow a = 0 \vee b = 0$

shows $a\text{-kernel } R \text{ rec-rng-of-frac rng-to-rng-of-frac} = \{0\}$

proof –

have $\{r \in \text{carrier } R. \text{rng-to-rng-of-frac } r = 0_{\text{rec-rng-of-frac}}\} \subseteq \{0\}$

proof(*rule subsetI*)

fix x

assume $a1: x \in \{r \in \text{carrier } R. \text{rng-to-rng-of-frac } r = 0_{\text{rec-rng-of-frac}}\}$

then have $(x, 1) \text{.}=\text{rel } (0, 1)$

using *rng-to-rng-of-frac-def rec-rng-of-frac-def eq-class-to-rel*

by *simp*

then obtain t **where** $f1:t \in S$ **and** $f2:t \otimes (1 \otimes x \oplus 1 \otimes 0) = 0$

using *rel-def*

by *auto*

have $f3:x \in \text{carrier } R$

using *a1*

by *simp*

then have $f4:t \otimes x = 0$

using *l-one r-zero f2*

by (*simp add: a-minus-def*)

have $t \neq 0$

using *f1 assms(1)*

by *auto*

then have $x = 0$

using *assms(2) f1 f3 f4 subset rev-subsetD*

by *auto*

thus $x \in \{0\}$

by *simp*

```

qed
have  $\{0\} \subseteq \{r \in \text{carrier } R. \text{rng-to-rng-of-frac } r = \mathbf{0}_{\text{rec-rng-of-frac}}\}$ 
  using subsetI rng-to-rng-of-frac-def rec-rng-of-frac-def
  by simp
then have  $\{r \in \text{carrier } R. \text{rng-to-rng-of-frac } r = \mathbf{0}_{\text{rec-rng-of-frac}}\} = \{0\}$ 
  using  $\langle \{r \in \text{carrier } R. \text{rng-to-rng-of-frac } r = \mathbf{0}_{\text{rec-rng-of-frac}}\} \subseteq \{0\} \rangle$ 
  by auto
thus ?thesis
  by (simp add: a-kernel-def kernel-def)
qed

end

end

```

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References

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