

Kraus Maps*

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Abstract

We formalize Kraus maps [1, Section 3], i.e., quantum channels of the form $\rho \mapsto \sum_x M_x \rho M_x^\dagger$ for suitable families of “Kraus operators” M_x . (In the finite-dimensional setting and the setting of separable Hilbert spaces, those are known to be equivalent to completely-positive maps, another common formalization of quantum channels.) Our results hold for arbitrary (i.e., not necessarily finite-dimensional or separable) Hilbert spaces.

Specifically, in theory `Kraus_Families`, we formalize the type (α, β, ξ) `kraus_family` of families of Kraus operators M_x (Kraus families for short), from trace-class operators on Hilbert space α to those on β , indexed by x of type ξ . This induces both a Kraus map $\rho \mapsto \sum_x M_x \rho M_x^\dagger$, as well as a quantum measurement with outcomes of type ξ .

We define and study various special Kraus families such as zero, identity, application of an operator, sum of two Kraus maps, sequential composition, infinite sum, random sampling, trace, tensor products, and complete measurements.

Furthermore, since working with explicit Kraus families can be cumbersome when the specific family of operators is not relevant, in theory `Kraus_Maps`, we define a Kraus map to be a function between two spaces that is of the form $\rho \mapsto \sum_x M_x \rho M_x^\dagger$ for some Kraus family, and restate our results in terms of such functions.

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1 Miscellaneous missing theorems

theory *Misc-Kraus-Maps*

imports

Hilbert-Space-Tensor-Product.Hilbert-Space-Tensor-Product

Hilbert-Space-Tensor-Product.Von-Neumann-Algebras

begin

unbundle *cblinfun-syntax*

lemma *abs-summable-norm*:

assumes $\langle f \text{ abs-summable-on } A \rangle$

shows $\langle (\lambda x. \text{norm } (f x)) \text{ abs-summable-on } A \rangle$

$\langle \text{proof} \rangle$

lemma *abs-summable-on-add*:

assumes $\langle f \text{ abs-summable-on } A \rangle$ **and** $\langle g \text{ abs-summable-on } A \rangle$

shows $\langle (\lambda x. f x + g x) \text{ abs-summable-on } A \rangle$

$\langle \text{proof} \rangle$

lemma *bdd-above-transform-mono-pos*:

assumes *bdd*: $\langle \text{bdd-above } ((\lambda x. g x) \text{ ' } M) \rangle$

assumes *gpos*: $\langle \bigwedge x. x \in M \implies g x \geq 0 \rangle$

assumes *mono*: $\langle \text{mono-on } (\text{Collect } ((\leq) 0)) f \rangle$

shows $\langle \text{bdd-above } ((\lambda x. f (g x)) \text{ ' } M) \rangle$

<proof>

lemma *Ex-iffI*:

assumes $\langle \bigwedge x. P\ x \implies Q\ (f\ x) \rangle$

assumes $\langle \bigwedge x. Q\ x \implies P\ (g\ x) \rangle$

shows $\langle Ex\ P \longleftrightarrow Ex\ Q \rangle$

<proof>

lemma *has-sum-Sigma'-banach*:

fixes $A :: 'a\ set$ **and** $B :: 'a \Rightarrow 'b\ set$

and $f :: 'a \Rightarrow 'b \Rightarrow 'c::banach$

assumes $\langle (\lambda(x,y). f\ x\ y)\ has-sum\ S \rangle$ $(Sigma\ A\ B)$

shows $\langle (\lambda x. infsum\ (f\ x)\ (B\ x))\ has-sum\ S \rangle\ A$

<proof>

lemma *infsum-Sigma-topological-monoid*:

fixes $A :: 'a\ set$ **and** $B :: 'a \Rightarrow 'b\ set$

and $f :: 'a \times 'b \Rightarrow 'c::\{topological-comm-monoid-add, t3-space\}$

assumes *summableAB*: f *summable-on* $(Sigma\ A\ B)$

assumes *summableB*: $\langle \bigwedge x. x \in A \implies (\lambda y. f\ (x, y))\ summable-on\ (B\ x) \rangle$

shows $infsum\ f\ (Sigma\ A\ B) = (\sum_{\infty x \in A}. \sum_{\infty y \in B\ x}. f\ (x, y))$

<proof>

lemma *flip-eq-const*: $\langle (\lambda y. y = x) = ((=)\ x) \rangle$

<proof>

lemma *sgn-ket[simp]*: $\langle sgn\ (ket\ x) = ket\ x \rangle$

<proof>

lemma *tensor-op-in-tensor-vn*:

assumes $\langle a \in A \rangle$ **and** $\langle b \in B \rangle$

shows $\langle a \otimes_o b \in A \otimes_{vN}\ B \rangle$

<proof>

lemma *commutant-tensor-vn-subset*:

assumes $\langle von-neumann-algebra\ A \rangle$ **and** $\langle von-neumann-algebra\ B \rangle$

shows $\langle commutant\ A \otimes_{vN}\ commutant\ B \subseteq commutant\ (A \otimes_{vN}\ B) \rangle$

<proof>

lemma *commutant-span[simp]*: $\langle commutant\ (span\ X) = commutant\ X \rangle$

<proof>

lemma *explicit-cblinfun-exists-0[simp]*: $\langle explicit-cblinfun-exists\ (\lambda- . 0) \rangle$

<proof>

lemma *explicit-cblinfun-0[simp]*: $\langle explicit-cblinfun\ (\lambda- . 0) = 0 \rangle$

<proof>

lemma *cnj-of-bool[simp]*: $\langle \text{cnj } (\text{of-bool } b) = \text{of-bool } b \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-single*:
fixes $f :: \langle - \Rightarrow - :: \{ \text{comm-monoid-add}, \text{t2-space} \} \rangle$
assumes $\langle \bigwedge j. j \neq i \implies j \in A \implies f j = 0 \rangle$
assumes $\langle s = (\text{if } i \in A \text{ then } f i \text{ else } 0) \rangle$
shows *HAS-SUM* $f A s$
 $\langle \text{proof} \rangle$

lemma *classical-operator-None[simp]*: $\langle \text{classical-operator } (\lambda -. \text{None}) = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-in-in-closedsubspace*:
assumes $\langle \text{has-sum-in } T f A l \rangle$
assumes $\langle \bigwedge x. x \in A \implies f x \in S \rangle$
assumes $\langle \text{closedin } T S \rangle$
assumes $\langle \text{csubspace } S \rangle$
shows $\langle l \in S \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-coordinatewise*:
 $\langle (f \text{ has-sum } s) A \longleftrightarrow (\forall i. ((\lambda x. f x i) \text{ has-sum } s i) A) \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-butterfly*:
 $\langle \text{butterfly } g h = (\text{one-dim-iso } g * \text{cnj } (\text{one-dim-iso } h)) *_C 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-tc-butterfly*:
fixes $g :: \langle 'a :: \text{one-dim} \rangle$ **and** $h :: \langle 'b :: \text{one-dim} \rangle$
shows $\langle \text{tc-butterfly } g h = (\text{one-dim-iso } g * \text{cnj } (\text{one-dim-iso } h)) *_C 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *one-dim-iso-of-real[simp]*: $\langle \text{one-dim-iso } (\text{of-real } x) = \text{of-real } x \rangle$
 $\langle \text{proof} \rangle$

lemma *filter-insert-if*:
 $\langle \text{Set.filter } P (\text{insert } x M) = (\text{if } P x \text{ then } \text{insert } x (\text{Set.filter } P M) \text{ else } \text{Set.filter } P M) \rangle$
 $\langle \text{proof} \rangle$

lemma *filter-empty[simp]*: $\langle \text{Set.filter } P \{\} = \{\} \rangle$
 $\langle \text{proof} \rangle$

lemma *has-sum-in-cong-neutral*:

fixes $f\ g :: \langle 'a \Rightarrow 'b :: \text{comm-monoid-add} \rangle$
assumes $\langle \bigwedge x. x \in T - S \implies g\ x = 0 \rangle$
assumes $\langle \bigwedge x. x \in S - T \implies f\ x = 0 \rangle$
assumes $\langle \bigwedge x. x \in S \cap T \implies f\ x = g\ x \rangle$
shows $\text{has-sum-in } X\ f\ S\ x \longleftrightarrow \text{has-sum-in } X\ g\ T\ x$
 $\langle \text{proof} \rangle$

lemma *infsun-in-cong-neutral*:
fixes $f\ g :: \langle 'a \Rightarrow 'b :: \text{comm-monoid-add} \rangle$
assumes $\langle \bigwedge x. x \in T - S \implies g\ x = 0 \rangle$
assumes $\langle \bigwedge x. x \in S - T \implies f\ x = 0 \rangle$
assumes $\langle \bigwedge x. x \in S \cap T \implies f\ x = g\ x \rangle$
shows $\langle \text{infsun-in } X\ f\ S = \text{infsun-in } X\ g\ T \rangle$
 $\langle \text{proof} \rangle$

lemma *filter-image*: $\langle \text{Set.filter } P\ (f\ ' X) = f\ ' (\text{Set.filter } (\lambda x. P\ (f\ x))\ X) \rangle$
 $\langle \text{proof} \rangle$

lemma *Sigma-image-left*: $\langle (\text{SIGMA } x:f'A. B\ x) = (\lambda(x,y). (f\ x, y))\ ' (\text{SIGMA } x:A. B\ (f\ x)) \rangle$
 $\langle \text{proof} \rangle$

lemma *finite-subset-filter-image*:
assumes *finite* B
assumes $\langle B \subseteq \text{Set.filter } P\ (f\ ' A) \rangle$
shows $\exists C \subseteq A. \text{finite } C \wedge B = f\ ' C$
 $\langle \text{proof} \rangle$

definition $\langle \text{card-le-1 } M \longleftrightarrow (\exists x. M \subseteq \{x\}) \rangle$

lemma *card-le-1-empty*[*iff*]: $\langle \text{card-le-1 } \{\} \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-singleton*[*iff*]: $\langle \text{card-le-1 } \{x\} \rangle$
 $\langle \text{proof} \rangle$

lemma *sgn-tensor-ell2*: $\langle \text{sgn } (h \otimes_s k) = \text{sgn } h \otimes_s \text{sgn } k \rangle$
 $\langle \text{proof} \rangle$

lemma *is-ortho-set-tensor*:
assumes $\langle \text{is-ortho-set } B \rangle$
assumes $\langle \text{is-ortho-set } C \rangle$
shows $\langle \text{is-ortho-set } ((\lambda(x, y). x \otimes_s y)\ ' (B \times C)) \rangle$
 $\langle \text{proof} \rangle$

end

2 Kraus families

theory *Kraus-Families*

imports

Wlog.Wlog

Hilbert-Space-Tensor-Product.Partial-Trace

Misc-Kraus-Maps

abbrevs

$=_{kr} = =_{kr}$ **and** $=_{kr} = =_{kr}$ **and** $*_{kr} = *_{kr}$

begin

unbundle *cblinfun-syntax*

2.1 Kraus families

definition $\langle kraus-family \mathfrak{E} \longleftrightarrow bdd-above ((\lambda F. \sum (E,x) \in F. E* \ o_{CL} \ E) \ ' \ {F. \ finite \ F \ \wedge \ F \subseteq \mathfrak{E}}) \ \wedge \ 0 \notin \text{fst} \ ' \ \mathfrak{E} \rangle$

for $\mathfrak{E} :: \langle (\text{--} :: \text{chilbert-space} \Rightarrow_{CL} \text{--} :: \text{chilbert-space}) \times \text{--} \rangle \text{ set}$

typedef (**overloaded**) $('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family} =$

$\langle \text{Collect kraus-family} :: (('a \Rightarrow_{CL} 'b) \times 'x) \text{ set set} \rangle$

$\langle \text{proof} \rangle$

setup-lifting *type-definition-kraus-family*

lemma *kraus-familyI:*

assumes $\langle bdd-above ((\lambda F. \sum (E,x) \in F. E* \ o_{CL} \ E) \ ' \ {F. \ finite \ F \ \wedge \ F \subseteq \mathfrak{E}}) \rangle$

assumes $\langle 0 \notin \text{fst} \ ' \ \mathfrak{E} \rangle$

shows $\langle kraus-family \ \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lift-definition *kf-apply* $:: \langle ('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family} \Rightarrow ('a, 'a) \text{ trace-class} \Rightarrow ('b, 'b) \text{ trace-class} \rangle$ **is**

$\langle \lambda \mathfrak{E} \ \varrho. (\sum_{\infty} E \in \mathfrak{E}. \text{sandwich-tc} (\text{fst} \ E) \ \varrho) \rangle$ $\langle \text{proof} \rangle$

notation *kf-apply* (**infixr** $\langle *_{kr} \rangle$ 70)

lemma *kraus-family-if-finite* [*iff*]: $\langle kraus-family \ \mathfrak{E} \rangle$ **if** $\langle \text{finite} \ \mathfrak{E} \rangle$ **and** $\langle 0 \notin \text{fst} \ ' \ \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-scaleC:*

shows $\langle kf-apply \ \mathfrak{E} (c *_{C} \ x) = c *_{C} \ kf-apply \ \mathfrak{E} \ x \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-abs-summable:*

assumes $\langle kraus-family \ \mathfrak{E} \rangle$

shows $\langle (\lambda (E,x). \text{sandwich-tc} \ E \ \varrho) \text{ abs-summable-on} \ \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-summable*:

shows $\langle (\lambda(E,x). \text{ sandwich-tc } E \ \varrho) \text{ summable-on } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
<proof>

lemma *kf-apply-has-sum*:

shows $\langle ((\lambda(E,x). \text{ sandwich-tc } E \ \varrho) \text{ has-sum } \text{kf-apply } \mathfrak{E} \ \varrho) (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
<proof>

lemma *kf-apply-plus-right*:

shows $\langle \text{kf-apply } \mathfrak{E} (x + y) = \text{kf-apply } \mathfrak{E} x + \text{kf-apply } \mathfrak{E} y \rangle$
<proof>

lemma *kf-apply-uminus-right*:

shows $\langle \text{kf-apply } \mathfrak{E} (-x) = - \text{kf-apply } \mathfrak{E} x \rangle$
<proof>

lemma *kf-apply-minus-right*:

shows $\langle \text{kf-apply } \mathfrak{E} (x - y) = \text{kf-apply } \mathfrak{E} x - \text{kf-apply } \mathfrak{E} y \rangle$
<proof>

lemma *kf-apply-pos*:

assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kf-apply } \mathfrak{E} \ \varrho \geq 0 \rangle$
<proof>

lemma *kf-apply-mono-right*:

assumes $\langle \varrho \geq \tau \rangle$
shows $\langle \text{kf-apply } \mathfrak{E} \ \varrho \geq \text{kf-apply } \mathfrak{E} \ \tau \rangle$
<proof>

lemma *kf-apply-geq-sum*:

assumes $\langle \varrho \geq 0 \rangle$ **and** $\langle M \subseteq \text{Rep-kraus-family } \mathfrak{E} \rangle$
shows $\langle \text{kf-apply } \mathfrak{E} \ \varrho \geq (\sum (E,-) \in M. \text{ sandwich-tc } E \ \varrho) \rangle$
<proof>

lift-definition *kf-domain* :: $\langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \Rightarrow 'x \text{ set} \rangle$ **is**
 $\langle \lambda \mathfrak{E}. \text{ snd } ' \mathfrak{E} \rangle$ *<proof>*

lemma *kf-apply-clinear*[iff]: $\langle \text{clinear } (\text{kf-apply } \mathfrak{E}) \rangle$

<proof>

lemma *kf-apply-0-right*[iff]: $\langle \text{kf-apply } \mathfrak{E} \ 0 = 0 \rangle$

<proof>

lift-definition *kf-operators* :: $\langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \Rightarrow ('a \Rightarrow_{CL} 'b) \text{ set} \rangle$ **is**
 $\langle \text{image fst} :: ('a \Rightarrow_{CL} 'b \times 'x) \text{ set} \Rightarrow ('a \Rightarrow_{CL} 'b) \text{ set} \rangle$ *proof*

2.2 Bound and norm

lift-definition *kf-bound* :: $\langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \Rightarrow ('a \Rightarrow_{CL} 'b) \rangle$ **is**
 $\langle \lambda \mathfrak{E}. \text{infsun-in cweak-operator-topology } (\lambda(E,x). E^* o_{CL} E) \mathfrak{E} \rangle$ *proof*

lemma *kf-bound-def'*:

$\langle \text{kf-bound } \mathfrak{E} = \text{Rep-cblinfun-wot } (\sum_{\infty} (E,x) \in \text{Rep-kraus-family } \mathfrak{E}. \text{compose-wot } (\text{adj-wot } (\text{Abs-cblinfun-wot } E)) (\text{Abs-cblinfun-wot } E)) \rangle$
proof

definition $\langle \text{kf-norm } \mathfrak{E} = \text{norm } (\text{kf-bound } \mathfrak{E}) \rangle$

lemma *kf-norm-sum-bdd*: $\langle \text{bdd-above } ((\lambda F. \text{norm } (\sum (E,x) \in F. E^* o_{CL} E)) \text{ ' } \{F. F \subseteq \text{Rep-kraus-family } \mathfrak{E} \wedge \text{finite } F\}) \rangle$
proof

lemma *kf-norm-geq0* [iff]:

shows $\langle \text{kf-norm } \mathfrak{E} \geq 0 \rangle$
proof

lemma *kf-bound-has-sum*:

shows $\langle \text{has-sum-in cweak-operator-topology } (\lambda(E,x). E^* o_{CL} E) (\text{Rep-kraus-family } \mathfrak{E}) (\text{kf-bound } \mathfrak{E}) \rangle$
proof

lemma *kraus-family-iff-summable*:

$\langle \text{kraus-family } \mathfrak{E} \longleftrightarrow \text{summable-on-in cweak-operator-topology } (\lambda(E,x). E^* o_{CL} E) \mathfrak{E} \wedge 0 \notin \text{fst ' } \mathfrak{E} \rangle$
proof

lemma *kraus-family-iff-summable'*:

$\langle \text{kraus-family } \mathfrak{E} \longleftrightarrow (\lambda(E,x). \text{Abs-cblinfun-wot } (E^* o_{CL} E)) \text{ summable-on } \mathfrak{E} \wedge 0 \notin \text{fst ' } \mathfrak{E} \rangle$
proof

lemma *kf-bound-summable*:

shows $\langle \text{summable-on-in cweak-operator-topology } (\lambda(E,x). E^* o_{CL} E) (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
proof

lemma *kf-bound-has-sum'*:

shows $\langle ((\lambda(E,x). \text{compose-wot } (\text{adj-wot } (\text{Abs-cblinfun-wot } E)) (\text{Abs-cblinfun-wot } E)) \text{ has-sum } \text{Abs-cblinfun-wot } (\text{kf-bound } \mathfrak{E})) (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
proof

lemma *kf-bound-summable'*:

$\langle (\lambda(E,x). \text{compose-wot } (\text{adj-wot } (\text{Abs-cblinfun-wot } E)) (\text{Abs-cblinfun-wot } E)) \text{ summable-on } \text{Rep-kraus-family } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-is-Sup*:

shows $\langle \text{is-Sup } ((\lambda F. \sum (E,x) \in F. E^* \text{ o}_{CL} E) \text{ ' } \{F. \text{finite } F \wedge F \subseteq \text{Rep-kraus-family } \mathfrak{E}\}) \rangle$
 $\langle \text{kf-bound } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-leqI*:

assumes $\langle \bigwedge F. \text{finite } F \implies F \subseteq \text{Rep-kraus-family } \mathfrak{E} \implies (\sum (E,x) \in F. E^* \text{ o}_{CL} E) \leq B \rangle$
shows $\langle \text{kf-bound } \mathfrak{E} \leq B \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-pos[iff]*: $\langle \text{kf-bound } \mathfrak{E} \geq 0 \rangle$

$\langle \text{proof} \rangle$

lemma *not-not-singleton-kf-norm-0*:

fixes $\mathfrak{E} :: \langle 'a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x \rangle \text{ kraus-family}$
assumes $\langle \neg \text{class.not-singleton TYPE('a)} \rangle$
shows $\langle \text{kf-norm } \mathfrak{E} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-sum-leqI*:

assumes $\langle \bigwedge F. \text{finite } F \implies F \subseteq \text{Rep-kraus-family } \mathfrak{E} \implies \text{norm } (\sum (E,x) \in F. E^* \text{ o}_{CL} E) \leq B \rangle$
shows $\langle \text{kf-norm } \mathfrak{E} \leq B \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-geq-sum*:

assumes $\langle M \subseteq \text{Rep-kraus-family } \mathfrak{E} \rangle$
shows $\langle (\sum (E,-) \in M. E^* \text{ o}_{CL} E) \leq \text{kf-bound } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-geq-norm-sum*:

assumes $\langle M \subseteq \text{Rep-kraus-family } \mathfrak{E} \rangle$
shows $\langle \text{norm } (\sum (E,-) \in M. E^* \text{ o}_{CL} E) \leq \text{kf-norm } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-finite*: $\langle \text{kf-bound } \mathfrak{E} = (\sum (E,x) \in \text{Rep-kraus-family } \mathfrak{E}. E^* \text{ o}_{CL} E) \rangle$ **if** $\langle \text{finite } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-finite*: $\langle \text{kf-norm } \mathfrak{E} = \text{norm } (\sum (E,x) \in \text{Rep-kraus-family } \mathfrak{E}. E^* \text{ o}_{CL} E) \rangle$

if $\langle \text{finite } (\text{Rep-kraus-family } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-bounded-pos*:
assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{norm } (\text{kf-apply } \mathfrak{E} \varrho) \leq \text{kf-norm } \mathfrak{E} * \text{norm } \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-bounded*:
— We suspect the factor 4 is not needed here but don't know how to prove that.
 $\langle \text{norm } (\text{kf-apply } \mathfrak{E} \varrho) \leq 4 * \text{kf-norm } \mathfrak{E} * \text{norm } \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-bounded-clinear*[*bounded-clinear*]:
shows $\langle \text{bounded-clinear } (\text{kf-apply } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-from-map*: $\langle \psi \cdot_C \text{kf-bound } \mathfrak{E} \varphi = \text{trace-tc } (\mathfrak{E} *_{k_r} \text{tc-butterfly } \varphi \psi) \rangle$
 $\langle \text{proof} \rangle$

lemma *trace-from-kf-bound*: $\langle \text{trace-tc } (\mathfrak{E} *_{k_r} \varrho) = \text{trace-tc } (\text{compose-tcr } (\text{kf-bound } \mathfrak{E}) \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-selfadjoint*[*iff*]: $\langle \text{selfadjoint } (\text{kf-bound } \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-leq-kf-norm-id*:
shows $\langle \text{kf-bound } \mathfrak{E} \leq \text{kf-norm } \mathfrak{E} *_R \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

2.3 Basic Kraus families

lemma *kf-emptyset*[*iff*]: $\langle \text{kraus-family } \{\} \rangle$
 $\langle \text{proof} \rangle$

instantiation *kraus-family* :: $(\text{hilbert-space}, \text{hilbert-space}, \text{type}) \langle \text{zero} \rangle$ **begin**

lift-definition *zero-kraus-family* :: $\langle ('a, 'b, 'x) \text{kraus-family} \rangle$ **is** $\langle \{\} \rangle$

$\langle \text{proof} \rangle$

instance $\langle \text{proof} \rangle$

end

lemma *kf-apply-0*[*simp*]: $\langle \text{kf-apply } 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-0*[*simp*]: $\langle \text{kf-bound } 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-0*[*simp*]: $\langle \text{kf-norm } 0 = 0 \rangle$

⟨proof⟩

lift-definition $kf\text{-of-op} :: \langle 'a :: \text{chilbert-space} \Rightarrow_{CL} 'b :: \text{chilbert-space} \Rightarrow ('a, 'b, \text{unit}) \text{ kraus-family} \rangle$
is

⟨ $\lambda E :: 'a \Rightarrow_{CL} 'b. \text{if } E = 0 \text{ then } \{\} \text{ else } \{(E, ())\}$ ⟩
⟨proof⟩

lemma $kf\text{-of-op}0[simp]: \langle kf\text{-of-op } 0 = 0 \rangle$
⟨proof⟩

lemma $kf\text{-of-op-norm}[simp]: \langle kf\text{-norm } (kf\text{-of-op } E) = (\text{norm } E)^2 \rangle$
⟨proof⟩

lemma $kf\text{-operators-in-kf-of-op}[simp]: \langle kf\text{-operators } (kf\text{-of-op } U) = (\text{if } U = 0 \text{ then } \{\} \text{ else } \{U\}) \rangle$
⟨proof⟩

lemma $kf\text{-domain-of-op}[simp]: \langle kf\text{-domain } (kf\text{-of-op } A) = \{\} \rangle$ **if** $\langle A \neq 0 \rangle$
⟨proof⟩

definition $\langle kf\text{-id} = kf\text{-of-op } id\text{-cblinfun} \rangle$

lemma $kf\text{-domain-id}[simp]: \langle kf\text{-domain } (kf\text{-id} :: ('a :: \{\text{chilbert-space, not-singleton}\}, -, -) \text{ kraus-family}) = \{\} \rangle$
⟨proof⟩

lemma $kf\text{-of-op-id}[simp]: \langle kf\text{-of-op } id\text{-cblinfun} = kf\text{-id} \rangle$
⟨proof⟩

lemma $kf\text{-norm-id-leq1}: \langle kf\text{-norm } kf\text{-id} \leq 1 \rangle$
⟨proof⟩

lemma $kf\text{-norm-id-eq1}[simp]: \langle kf\text{-norm } (kf\text{-id} :: ('a :: \{\text{chilbert-space, not-singleton}\}, 'a, \text{unit}) \text{ kraus-family}) = 1 \rangle$
⟨proof⟩

lemma $kf\text{-operators-in-kf-id}[simp]: \langle kf\text{-operators } kf\text{-id} = (\text{if } (id\text{-cblinfun} :: 'a :: \text{chilbert-space} \Rightarrow_{CL} -) = 0 \text{ then } \{\} \text{ else } \{id\text{-cblinfun} :: 'a \Rightarrow_{CL} -\}) \rangle$
⟨proof⟩

instantiation $\text{kraus-family} :: (\text{chilbert-space}, \text{chilbert-space}, \text{type}) \text{ scaleR } \text{begin}$

lift-definition $\text{scaleR-kraus-family} :: \langle \text{real} \Rightarrow ('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family} \Rightarrow ('a, 'b, 'x) \text{ kraus-family} \rangle$ **is**

⟨ $\lambda r \in \mathfrak{C}. \text{if } r \leq 0 \text{ then } \{\} \text{ else } (\lambda(E, x). (\text{sqrt } r *_{\mathbb{R}} E, x)) \text{ ' } \mathfrak{C} \rangle$
⟨proof⟩

instance ⟨proof⟩

end

lemma $kf\text{-scale-apply}$:

assumes $\langle r \geq 0 \rangle$
shows $\langle \text{kf-apply } (r *_R \mathfrak{E}) \varrho = r *_R \text{kf-apply } \mathfrak{E} \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-scale[simp]*: $\langle \text{kf-bound } (c *_R \mathfrak{E}) = c *_R \text{kf-bound } \mathfrak{E} \rangle$ **if** $\langle c \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-scale[simp]*: $\langle \text{kf-norm } (c *_R \mathfrak{E}) = c * \text{kf-norm } \mathfrak{E} \rangle$ **if** $\langle c \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-of-op-apply*: $\langle \text{kf-apply } (\text{kf-of-op } E) \varrho = \text{sandwich-tc } E \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-id-apply[simp]*: $\langle \text{kf-apply } \text{kf-id } \varrho = \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-scale-apply-neg*:
assumes $\langle r \leq 0 \rangle$
shows $\langle \text{kf-apply } (r *_R \mathfrak{E}) = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-0-left[iff]*: $\langle \text{kf-apply } 0 \varrho = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-of-op[simp]*: $\langle \text{kf-bound } (\text{kf-of-op } A) = A * o_{CL} A \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-id[simp]*: $\langle \text{kf-bound } \text{kf-id} = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

2.4 Filtering

lift-definition *kf-filter* :: $\langle ('x \Rightarrow \text{bool}) \Rightarrow ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'x) \text{kraus-family} \Rightarrow ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'x) \text{kraus-family} \rangle$ **is**
 $\langle \lambda(P::'x \Rightarrow \text{bool}) \mathfrak{E}. \text{Set.filter } (\lambda(E,x). P x) \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-false[simp]*: $\langle \text{kf-filter } (\lambda-. \text{False}) \mathfrak{E} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-true[simp]*: $\langle \text{kf-filter } (\lambda-. \text{True}) \mathfrak{E} = \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

definition $\langle \text{kf-apply-on } \mathfrak{E} S = \text{kf-apply } (\text{kf-filter } (\lambda x. x \in S)) \mathfrak{E} \rangle$

notation *kf-apply-on* ($- *_k r @-/ - [71, 1000, 70] 70$)

lemma *kf-apply-on-pos*:
assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kf-apply-on } \mathfrak{E} X \varrho \geq 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-bounded-clinear*[*bounded-clinear*]:

shows $\langle \text{bounded-clinear } (kf\text{-apply-on } \mathfrak{E} X) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-filter-twice*:

$\langle kf\text{-filter } P (kf\text{-filter } Q \mathfrak{E}) = kf\text{-filter } (\lambda x. P x \wedge Q x) \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-union-has-sum*:

assumes $\langle \bigwedge X Y. X \in F \implies Y \in F \implies X \neq Y \implies \text{disjnt } X Y \rangle$

shows $\langle ((\lambda X. kf\text{-apply-on } \mathfrak{E} X \varrho) \text{ has-sum } (kf\text{-apply-on } \mathfrak{E} (\bigcup F) \varrho)) F \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-empty*[*simp*]: $\langle kf\text{-apply-on } E \{\} \varrho = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-union-eqI*:

assumes $\langle \bigwedge X Y. (X, Y) \in F \implies kf\text{-apply-on } \mathfrak{E} X \varrho = kf\text{-apply-on } \mathfrak{F} Y \varrho \rangle$

assumes $\langle \bigwedge X Y X' Y'. (X, Y) \in F \implies (X', Y') \in F \implies (X, Y) \neq (X', Y') \implies \text{disjnt } X X' \rangle$

assumes $\langle \bigwedge X Y X' Y'. (X, Y) \in F \implies (X', Y') \in F \implies (X, Y) \neq (X', Y') \implies \text{disjnt } Y Y' \rangle$

assumes $XX: \langle XX = \bigcup (\text{fst } ' F) \rangle$ **and** $YY: \langle YY = \bigcup (\text{snd } ' F) \rangle$

shows $\langle kf\text{-apply-on } \mathfrak{E} XX \varrho = kf\text{-apply-on } \mathfrak{F} YY \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-UNIV*[*simp*]: $\langle kf\text{-apply-on } \mathfrak{E} UNIV = kf\text{-apply } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-CARD-1*[*simp*]: $\langle (\lambda \mathfrak{E}. kf\text{-apply-on } \mathfrak{E} \{x::: CARD-1\}) = kf\text{-apply} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-union-summable-on*:

assumes $\langle \bigwedge X Y. X \in F \implies Y \in F \implies X \neq Y \implies \text{disjnt } X Y \rangle$

shows $\langle (\lambda X. kf\text{-apply-on } \mathfrak{E} X \varrho) \text{ summable-on } F \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-union-infsum*:

assumes $\langle \bigwedge X Y. X \in F \implies Y \in F \implies X \neq Y \implies \text{disjnt } X Y \rangle$

shows $\langle (\sum_{\infty} X \in F. kf\text{-apply-on } \mathfrak{E} X \varrho) = kf\text{-apply-on } \mathfrak{E} (\bigcup F) \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *kf-bound-filter*:

$\langle kf\text{-bound } (kf\text{-filter } P \mathfrak{E}) \leq kf\text{-bound } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-filter*:

$\langle \text{kf-norm } (\text{kf-filter } P \mathfrak{E}) \leq \text{kf-norm } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-domain-filter[simp]*:

$\langle \text{kf-domain } (\text{kf-filter } P E) = \text{Set.filter } P (\text{kf-domain } E) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-0-right[simp]*: $\langle \text{kf-filter } P 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-0-right[simp]*: $\langle \text{kf-apply-on } \mathfrak{E} X 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-0-left[simp]*: $\langle \text{kf-apply-on } 0 X \varrho = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-mono3*:

assumes $\langle \varrho \leq \sigma \rangle$

shows $\langle \mathfrak{E} *_{kr} @X \varrho \leq \mathfrak{E} *_{kr} @X \sigma \rangle$

$\langle \text{proof} \rangle$

lemma *kf-apply-on-mono2*:

assumes $\langle X \subseteq Y \rangle$ **and** $\langle \varrho \geq 0 \rangle$

shows $\langle \mathfrak{E} *_{kr} @X \varrho \leq \mathfrak{E} *_{kr} @Y \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *kf-operators-filter*: $\langle \text{kf-operators } (\text{kf-filter } P \mathfrak{E}) \subseteq \text{kf-operators } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-equal-if-filter-equal*:

assumes $\langle \bigwedge x. \text{kf-filter } ((=)x) \mathfrak{E} = \text{kf-filter } ((=)x) \mathfrak{F} \rangle$

shows $\langle \mathfrak{E} = \mathfrak{F} \rangle$

$\langle \text{proof} \rangle$

2.5 Equivalence

definition $\langle \text{kf-eq-weak } \mathfrak{E} \mathfrak{F} \longleftrightarrow \text{kf-apply } \mathfrak{E} = \text{kf-apply } \mathfrak{F} \rangle$

definition $\langle \text{kf-eq } \mathfrak{E} \mathfrak{F} \longleftrightarrow (\forall x. \text{kf-apply-on } \mathfrak{E} \{x\} = \text{kf-apply-on } \mathfrak{F} \{x\}) \rangle$

open-bundle *kraus-map-syntax* **begin**

notation *kf-eq-weak* (**infix** $=_{kr}$ 50)

notation *kf-eq* (**infix** \equiv_{kr} 50)

notation *kf-apply* (**infixr** $\langle *_{kr} \rangle$ 70)

notation *kf-apply-on* ($- *_{kr} @- / - [71, 1000, 70] 70$)

end

lemma *kf-eq-weak-refl*[*iff*]: $\langle x =_{kr} x \rangle$
<proof>

lemma *kf-eq-weak-refl0*[*iff*]: $\langle 0 =_{kr} 0 \rangle$
<proof>

lemma *kf-bound-cong*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \text{kf-bound } \mathfrak{E} = \text{kf-bound } \mathfrak{F} \rangle$
<proof>

lemma *kf-norm-cong*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \text{kf-norm } \mathfrak{E} = \text{kf-norm } \mathfrak{F} \rangle$
<proof>

lemma *kf-eq-weakI*:
assumes $\langle \bigwedge \varrho. \varrho \geq 0 \implies \mathfrak{E} *_{kr} \varrho = \mathfrak{F} *_{kr} \varrho \rangle$
shows $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
<proof>

lemma *kf-eqI*:
assumes $\langle \bigwedge x \varrho. \varrho \geq 0 \implies \mathfrak{E} *_{kr} @\{x\} \varrho = \mathfrak{F} *_{kr} @\{x\} \varrho \rangle$
shows $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
<proof>

lemma *kf-eq-weak-trans*[*trans*]:
 $\langle F =_{kr} G \implies G =_{kr} H \implies F =_{kr} H \rangle$
<proof>

lemma *kf-apply-eqI*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} *_{kr} \varrho = \mathfrak{F} *_{kr} \varrho \rangle$
<proof>

lemma *kf-eq-imp-eq-weak*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
<proof>

lemma *kf-filter-cong-eq*[*cong*]:
assumes $\langle \mathfrak{E} = \mathfrak{F} \rangle$
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{E} \implies P x = Q x \rangle$
shows $\langle \text{kf-filter } P \mathfrak{E} = \text{kf-filter } Q \mathfrak{F} \rangle$
<proof>

lemma *kf-filter-cong*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{E} \implies P x = Q x \rangle$
shows $\langle \text{kf-filter } P \ \mathfrak{E} \equiv_{kr} \text{kf-filter } Q \ \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-cong-weak*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{E} \implies P x = Q x \rangle$
shows $\langle \text{kf-filter } P \ \mathfrak{E} =_{kr} \text{kf-filter } Q \ \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-eq-refl[iff]*: $\langle \mathfrak{E} \equiv_{kr} \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-eq-trans[trans]*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
assumes $\langle \mathfrak{F} \equiv_{kr} \mathfrak{G} \rangle$
shows $\langle \mathfrak{E} \equiv_{kr} \mathfrak{G} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-eq-sym[sym]*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{F} \equiv_{kr} \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-eq-weak-imp-eq-CARD-1*:
fixes $\mathfrak{E} \ \mathfrak{F} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x::\text{CARD-1}) \text{ kraus-family} \rangle$
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-on-eqI-filter*:
assumes $\langle \text{kf-filter } (\lambda x. x \in X) \ \mathfrak{E} \equiv_{kr} \text{kf-filter } (\lambda x. x \in X) \ \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} *_{kr} @X \ \varrho = \mathfrak{F} *_{kr} @X \ \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-on-eqI*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{E} *_{kr} @X \ \varrho = \mathfrak{F} *_{kr} @X \ \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-eq0I*:
assumes $\langle \mathfrak{E} =_{kr} 0 \rangle$
shows $\langle \mathfrak{E} *_{kr} \varrho = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-eq-weak0-imp-kf-eq0*:

assumes $\langle \mathfrak{E} =_{kr} 0 \rangle$
shows $\langle \mathfrak{E} \equiv_{kr} 0 \rangle$
 $\langle proof \rangle$

lemma *kf-apply-on-eq0I*:
assumes $\langle \mathfrak{E} =_{kr} 0 \rangle$
shows $\langle \mathfrak{E} *_{kr} @X \varrho = 0 \rangle$
 $\langle proof \rangle$

lemma *kf-filter-to-domain[simp]*:
 $\langle kf-filter (\lambda x. x \in kf-domain \mathfrak{E}) \mathfrak{E} = \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *kf-eq-0-iff-eq-0*: $\langle E =_{kr} 0 \longleftrightarrow E = 0 \rangle$
 $\langle proof \rangle$

lemma *in-kf-domain-iff-apply-nonzero*:
 $\langle x \in kf-domain \mathfrak{E} \longleftrightarrow kf-apply-on \mathfrak{E} \{x\} \neq 0 \rangle$
 $\langle proof \rangle$

lemma *kf-domain-cong*:
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
shows $\langle kf-domain \mathfrak{E} = kf-domain \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-eq-weak-sym[sym]*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle \mathfrak{F} =_{kr} \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *kf-eqI-from-filter-eq-weak*:
assumes $\langle \bigwedge x. kf-filter ((=) x) E =_{kr} kf-filter ((=) x) F \rangle$
shows $\langle E \equiv_{kr} F \rangle$
 $\langle proof \rangle$

lemma *kf-eq-weak-from-separatingI*:
fixes $E :: \langle ('q :: chilbert-space, 'r :: chilbert-space, 'x) kraus-family \rangle$
and $F :: \langle ('q, 'r, 'y) kraus-family \rangle$
assumes $\langle separating-set (bounded-clinear :: (('q, 'q) trace-class \Rightarrow ('r, 'r) trace-class) \Rightarrow bool) \rangle$
 S
assumes $\langle \bigwedge \varrho. \varrho \in S \Longrightarrow E *_{kr} \varrho = F *_{kr} \varrho \rangle$
shows $\langle E \equiv_{kr} F \rangle$
 $\langle proof \rangle$

lemma *kf-eq-weak-eq-trans[trans]*: $\langle a =_{kr} b \Longrightarrow b \equiv_{kr} c \Longrightarrow a =_{kr} c \rangle$

<proof>

lemma *kf-eq-eq-weak-trans*[*trans*]: $\langle a \equiv_{kr} b \implies b =_{kr} c \implies a =_{kr} c \rangle$
<proof>

instantiation *kraus-family* :: (*chilbert-space, chilbert-space, type*) *preorder begin*

definition *less-eq-kraus-family* **where** $\langle \mathfrak{E} \leq \mathfrak{F} \longleftrightarrow (\forall x. \forall \rho \geq 0. \text{kf-apply-on } \mathfrak{E} \{x\} \rho \leq \text{kf-apply-on } \mathfrak{F} \{x\} \rho) \rangle$

definition *less-kraus-family* **where** $\langle \mathfrak{E} < \mathfrak{F} \longleftrightarrow \mathfrak{E} \leq \mathfrak{F} \wedge \neg \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$

lemma *kf-antisym*: $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \longleftrightarrow \mathfrak{E} \leq \mathfrak{F} \wedge \mathfrak{F} \leq \mathfrak{E} \rangle$

for $\mathfrak{E} \mathfrak{F} :: \langle ('a, 'b, 'c) \text{ kraus-family} \rangle$

<proof>

instance

<proof>

end

lemma *kf-apply-on-mono1*:

assumes $\langle \mathfrak{E} \leq \mathfrak{F} \rangle$ **and** $\langle \rho \geq 0 \rangle$

shows $\langle \mathfrak{E} *_{kr} @X \rho \leq \mathfrak{F} *_{kr} @X \rho \rangle$

<proof>

lemma *kf-apply-mono-left*: $\langle \mathfrak{E} \leq \mathfrak{F} \implies \rho \geq 0 \implies \mathfrak{E} *_{kr} \rho \leq \mathfrak{F} *_{kr} \rho \rangle$

<proof>

lemma *kf-apply-mono*:

assumes $\langle \rho \geq 0 \rangle$

assumes $\langle \mathfrak{E} \leq \mathfrak{F} \rangle$ **and** $\langle \rho \leq \sigma \rangle$

shows $\langle \mathfrak{E} *_{kr} \rho \leq \mathfrak{F} *_{kr} \sigma \rangle$

<proof>

lemma *kf-apply-on-mono*:

assumes $\langle \rho \geq 0 \rangle$

assumes $\langle \mathfrak{E} \leq \mathfrak{F} \rangle$ **and** $\langle X \subseteq Y \rangle$ **and** $\langle \rho \leq \sigma \rangle$

shows $\langle \mathfrak{E} *_{kr} @X \rho \leq \mathfrak{F} *_{kr} @Y \sigma \rangle$

<proof>

lemma *kf-one-dim-is-id*[*simp*]:

fixes $\mathfrak{E} :: \langle ('a::\text{one-dim}, 'a::\text{one-dim}, 'x) \text{ kraus-family} \rangle$

shows $\langle \mathfrak{E} =_{kr} \text{kf-norm } \mathfrak{E} *_{kr} \text{kf-id} \rangle$

<proof>

2.6 Mapping and flattening

definition *kf-similar-elements* :: $\langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \implies ('a \Rightarrow_{CL} 'b) \implies ('a \Rightarrow_{CL} 'b \times 'x) \text{ set} \rangle$ **where**

— All elements of the Kraus family that are equal up to rescaling (and belong to the same output)

$\langle \text{kf-similar-elements } \mathfrak{E} E = \{(F, x) \in \text{Rep-kraus-family } \mathfrak{E}. (\exists r > 0. E = r *_{kr} F)\} \rangle$

definition *kf-element-weight where*

— The total weight (norm of the square) of all similar elements

$\langle \text{kf-element-weight } \mathfrak{E} \ E = (\sum_{\infty} (F, -) \in \text{kf-similar-elements } \mathfrak{E} \ E. \text{ norm } (F^* \ o_{CL} \ F)) \rangle$

lemma *kf-element-weight-geq0[simp]*: $\langle \text{kf-element-weight } \mathfrak{E} \ E \geq 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-similar-elements-abs-summable*:

fixes $\mathfrak{E} :: \langle ('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family} \rangle$

shows $\langle (\lambda(F, -). F^* \ o_{CL} \ F) \text{ abs-summable-on } (\text{kf-similar-elements } \mathfrak{E} \ E) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-similar-elements-kf-operators*:

assumes $\langle (F, x) \in \text{kf-similar-elements } \mathfrak{E} \ E \rangle$

shows $\langle F \in \text{span } (\text{kf-operators } \mathfrak{E}) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-element-weight-neq0*: $\langle \text{kf-element-weight } \mathfrak{E} \ E \neq 0 \rangle$

if $\langle (E, x) \in \text{Rep-kraus-family } \mathfrak{E} \rangle$ **and** $\langle E \neq 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-element-weight-0-left[simp]*: $\langle \text{kf-element-weight } 0 \ E = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-element-weight-0-right[simp]*: $\langle \text{kf-element-weight } E \ 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kf-element-weight-scale*:

assumes $\langle r > 0 \rangle$

shows $\langle \text{kf-element-weight } \mathfrak{E} \ (r \ *_R \ E) = \text{kf-element-weight } \mathfrak{E} \ E \rangle$

$\langle \text{proof} \rangle$

lemma *kf-element-weight-kf-operators*:

assumes $\langle \text{kf-element-weight } \mathfrak{E} \ E \neq 0 \rangle$

shows $\langle E \in \text{span } (\text{kf-operators } \mathfrak{E}) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-map-aux*:

fixes $f :: \langle 'x \Rightarrow 'y \rangle$ **and** $\mathfrak{E} :: \langle ('a :: \text{chilbert-space}, 'b :: \text{chilbert-space}, 'x) \text{ kraus-family} \rangle$

defines $\langle B \equiv \text{kf-bound } \mathfrak{E} \rangle$

defines $\langle \text{filtered } y \equiv \text{kf-filter } (\lambda x. f \ x = y) \ \mathfrak{E} \rangle$

defines $\langle \text{flattened} \equiv \{ (E, y). \text{ norm } (E^* \ o_{CL} \ E) = \text{kf-element-weight } (\text{filtered } y) \ E \wedge E \neq 0 \} \rangle$

defines $\langle \text{good} \equiv (\lambda(E, y). (\text{norm } E)^2 = \text{kf-element-weight } (\text{filtered } y) \ E \wedge E \neq 0) \rangle$

shows $\langle \text{has-sum-in cweak-operator-topology } (\lambda(E, -). E^* \ o_{CL} \ E) \ \text{flattened } B \rangle$ **(is ?has-sum)**

and $\langle \text{snd } ' \ (\text{SIGMA } (E, y): \text{Collect good. kf-similar-elements } (\text{filtered } y) \ E)$

$= \{ (F, x). (F, x) \in \text{Rep-kraus-family } \mathfrak{E} \wedge F \neq 0 \} \rangle$ **(is ?snd-sigma)**

and $\langle \text{inj-on snd } (\text{SIGMA } p: \text{Collect good. kf-similar-elements } (\text{filtered } (\text{snd } p)) \ (\text{fst } p)) \rangle$ **(is**

?inj-snd)
 ⟨proof⟩

lift-definition $kf\text{-map} :: \langle 'x \Rightarrow 'y \rangle \Rightarrow \langle 'a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'x \rangle \text{ kraus-family} \Rightarrow \langle 'a, 'b, 'y \rangle \text{ kraus-family} \rangle$ **is**
 $\langle \lambda f \mathfrak{E}. \{(E, y). \text{norm } (E^* \circ_{CL} E) = \text{kf-element-weight } (\text{kf-filter } (\lambda x. f x = y) \mathfrak{E}) E \wedge E \neq 0\} \rangle$
 ⟨proof⟩

lemma

fixes $\mathfrak{E} :: \langle 'a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'x \rangle \text{ kraus-family} \rangle$
shows $kf\text{-apply-map}[simp]: \langle \text{kf-apply } (\text{kf-map } f \mathfrak{E}) = \text{kf-apply } \mathfrak{E} \rangle$
and $kf\text{-map-bound}: \langle \text{kf-bound } (\text{kf-map } f \mathfrak{E}) = \text{kf-bound } \mathfrak{E} \rangle$
and $kf\text{-map-norm}[simp]: \langle \text{kf-norm } (\text{kf-map } f \mathfrak{E}) = \text{kf-norm } \mathfrak{E} \rangle$
 ⟨proof⟩

abbreviation $\langle \text{kf-flatten} \equiv \text{kf-map } (\lambda-. ()) \rangle$

Like $kf\text{-map}$, but with a much simpler definition. However, only makes sense for injective functions.

lift-definition $kf\text{-map-inj} :: \langle 'x \Rightarrow 'y \rangle \Rightarrow \langle 'a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'x \rangle \text{ kraus-family} \Rightarrow \langle 'a, 'b, 'y \rangle \text{ kraus-family} \rangle$ **is**
 $\langle \lambda f \mathfrak{E}. (\lambda(E, x). (E, f x)) \text{ ' } \mathfrak{E} \rangle$
 ⟨proof⟩

lemma $kf\text{-element-weight-map-inj}$:

assumes $\langle \text{inj-on } f \text{ (kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-element-weight } (\text{kf-map-inj } f \mathfrak{E}) E = \text{kf-element-weight } \mathfrak{E} E \rangle$
 ⟨proof⟩

lemma $kf\text{-eq-weak-kf-map-left}$: $\langle \text{kf-map } f F =_{kr} G \rangle$ **if** $\langle F =_{kr} G \rangle$
 ⟨proof⟩

lemma $kf\text{-eq-weak-kf-map-right}$: $\langle F =_{kr} \text{kf-map } f G \rangle$ **if** $\langle F =_{kr} G \rangle$
 ⟨proof⟩

lemma $kf\text{-filter-map}$:

fixes $\mathfrak{E} :: \langle 'a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'x \rangle \text{ kraus-family} \rangle$
shows $\langle \text{kf-filter } P \text{ (kf-map } f \mathfrak{E}) = \text{kf-map } f \text{ (kf-filter } (\lambda x. P (f x)) \mathfrak{E}) \rangle$
 ⟨proof⟩

lemma $kf\text{-filter-map-inj}$:

fixes $\mathfrak{E} :: \langle 'a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'x \rangle \text{ kraus-family} \rangle$
shows $\langle \text{kf-filter } P \text{ (kf-map-inj } f \mathfrak{E}) = \text{kf-map-inj } f \text{ (kf-filter } (\lambda x. P (f x)) \mathfrak{E}) \rangle$
 ⟨proof⟩

lemma $kf\text{-map-kf-map-inj-comp}$:

assumes $\langle \text{inj-on } f \text{ (kf-domain } \mathfrak{E}) \rangle$

shows $\langle \text{kf-map } g \ (\text{kf-map-inj } f \ \mathfrak{E}) = \text{kf-map } (g \circ f) \ \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-element-weight-eqweak0*:
assumes $\langle \mathfrak{E} =_{k_r} 0 \rangle$
shows $\langle \text{kf-element-weight } \mathfrak{E} \ E = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-kf-map-comp*:
assumes $\langle \text{inj-on } g \ (f \ ' \ \text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-map-inj } g \ (\text{kf-map } f \ \mathfrak{E}) = \text{kf-map } (g \circ f) \ \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-map-inj[simp]*:
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-map-inj } f \ \mathfrak{E} \ *_{k_r} \varrho = \mathfrak{E} \ *_{k_r} \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-eq-kf-map*:
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-map-inj } f \ \mathfrak{E} \equiv_{k_r} \text{kf-map } f \ \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-id[simp]*: $\langle \text{kf-map-inj id } \mathfrak{E} = \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-id*: $\langle \text{kf-map id } \mathfrak{E} \equiv_{k_r} \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-bound[simp]*:
fixes $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \ \text{kraus-family} \rangle$
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-bound } (\text{kf-map-inj } f \ \mathfrak{E}) = \text{kf-bound } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-norm[simp]*:
fixes $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \ \text{kraus-family} \rangle$
assumes $\langle \text{inj-on } f \ (\text{kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-norm } (\text{kf-map-inj } f \ \mathfrak{E}) = \text{kf-norm } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-cong-weak*:
assumes $\langle \mathfrak{E} =_{k_r} \mathfrak{F} \rangle$
shows $\langle \text{kf-map } f \ \mathfrak{E} =_{k_r} \text{kf-map } g \ \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-flatten-cong-weak*:

assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle kf\text{-flatten } \mathfrak{E} =_{kr} kf\text{-flatten } \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-flatten-cong*:
assumes $\langle \mathfrak{E} =_{kr} \mathfrak{F} \rangle$
shows $\langle kf\text{-flatten } \mathfrak{E} \equiv_{kr} kf\text{-flatten } \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-map-twice*:
 $\langle kf\text{-map } f (kf\text{-map } g \mathfrak{E}) \equiv_{kr} kf\text{-map } (f \circ g) \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *kf-map-cong*:
assumes $\langle \bigwedge x. x \in kf\text{-domain } \mathfrak{E} \implies f x = g x \rangle$
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{F} \rangle$
shows $\langle kf\text{-map } f \mathfrak{E} \equiv_{kr} kf\text{-map } g \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-map-inj-cong-eq*:
assumes $\langle \bigwedge x. x \in kf\text{-domain } \mathfrak{E} \implies f x = g x \rangle$
assumes $\langle \mathfrak{E} = \mathfrak{F} \rangle$
shows $\langle kf\text{-map-inj } f \mathfrak{E} = kf\text{-map-inj } g \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-domain-map[simp]*:
 $\langle kf\text{-domain } (kf\text{-map } f \mathfrak{E}) = f \text{ ` } kf\text{-domain } \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *kf-apply-on-map[simp]*:
 $\langle (kf\text{-map } f E) *_{kr} @X \varrho = E *_{kr} @(f \text{ ` } X) \varrho \rangle$
 $\langle proof \rangle$

lemma *kf-apply-on-map-inj[simp]*:
assumes $\langle inj\text{-on } f ((f \text{ ` } X) \cap kf\text{-domain } E) \rangle$
shows $\langle kf\text{-map-inj } f E *_{kr} @X \varrho = E *_{kr} @(f \text{ ` } X) \varrho \rangle$
 $\langle proof \rangle$

lemma *kf-map0[simp]*: $\langle kf\text{-map } f 0 = 0 \rangle$
 $\langle proof \rangle$

lemma *kf-map-inj-kr-eq-weak*:
assumes $\langle inj\text{-on } f (kf\text{-domain } \mathfrak{E}) \rangle$
shows $\langle kf\text{-map-inj } f \mathfrak{E} =_{kr} \mathfrak{E} \rangle$

⟨proof⟩

lemma *kf-map-inj-0[simp]*: ⟨*kf-map-inj* *f* 0 = 0⟩
⟨proof⟩

lemma *kf-domain-map-inj[simp]*: ⟨*kf-domain* (*kf-map-inj* *f* \mathfrak{E}) = *f* ‘*kf-domain* \mathfrak{E} ⟩
⟨proof⟩

lemma *kf-operators-kf-map*:
⟨*kf-operators* (*kf-map* *f* \mathfrak{E}) \subseteq *span* (*kf-operators* \mathfrak{E})⟩
⟨proof⟩

lemma *kf-operators-kf-map-inj[simp]*: ⟨*kf-operators* (*kf-map-inj* *f* \mathfrak{E}) = *kf-operators* \mathfrak{E} ⟩
⟨proof⟩

2.7 Addition

lift-definition *kf-plus* :: ⟨('a::chilbert-space, 'b::chilbert-space, 'x) *kraus-family* \Rightarrow ('a,'b,'y) *kraus-family* \Rightarrow ('a,'b,'x+'y) *kraus-family*⟩ **is**
⟨ $\lambda \mathfrak{E} \mathfrak{F}. (\lambda (E,x). (E, \text{Inl } x)) \text{ ‘ } \mathfrak{E} \cup (\lambda (F,y). (F, \text{Inr } y)) \text{ ‘ } \mathfrak{F}$ ⟩
⟨proof⟩

instantiation *kraus-family* :: (*chilbert-space*, *chilbert-space*, *type*) **plus begin**

definition *plus-kraus-family* **where** $\langle \mathfrak{E} + \mathfrak{F} = \text{kf-map } (\lambda xy. \text{case } xy \text{ of } \text{Inl } x \Rightarrow x \mid \text{Inr } y \Rightarrow y)$
(*kf-plus* \mathfrak{E} \mathfrak{F})⟩

instance⟨proof⟩

end

lemma *kf-plus-apply*:
fixes $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('a, 'b, 'y) \text{ kraus-family} \rangle$
shows $\langle \text{kf-apply } (\text{kf-plus } \mathfrak{E} \mathfrak{F}) \varrho = \text{kf-apply } \mathfrak{E} \varrho + \text{kf-apply } \mathfrak{F} \varrho \rangle$
⟨proof⟩

lemma *kf-plus-apply'*: $\langle (\mathfrak{E} + \mathfrak{F}) *_{kr} \varrho = \mathfrak{E} *_{kr} \varrho + \mathfrak{F} *_{kr} \varrho \rangle$
⟨proof⟩

lemma *kf-plus-0-left[simp]*: $\langle \text{kf-plus } 0 \mathfrak{E} = \text{kf-map-inj } \text{Inr } \mathfrak{E} \rangle$
⟨proof⟩

lemma *kf-plus-0-right[simp]*: $\langle \text{kf-plus } \mathfrak{E} 0 = \text{kf-map-inj } \text{Inl } \mathfrak{E} \rangle$
⟨proof⟩

lemma *kf-plus-0-left'[simp]*: $\langle 0 + \mathfrak{E} \equiv_{kr} \mathfrak{E} \rangle$
⟨proof⟩

lemma *kf-plus-0-right'*: $\langle \mathfrak{E} + 0 \equiv_{kr} \mathfrak{E} \rangle$
⟨proof⟩

lemma *kf-plus-bound*: $\langle \text{kf-bound } (\text{kf-plus } \mathfrak{E} \ \mathfrak{F}) = \text{kf-bound } \mathfrak{E} + \text{kf-bound } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-plus-bound'*: $\langle \text{kf-bound } (\mathfrak{E} + \mathfrak{F}) = \text{kf-bound } \mathfrak{E} + \text{kf-bound } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-triangle*: $\langle \text{kf-norm } (\text{kf-plus } \mathfrak{E} \ \mathfrak{F}) \leq \text{kf-norm } \mathfrak{E} + \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-triangle'*: $\langle \text{kf-norm } (\mathfrak{E} + \mathfrak{F}) \leq \text{kf-norm } \mathfrak{E} + \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-plus-map-both*:
 $\langle \text{kf-plus } (\text{kf-map } f \ \mathfrak{E}) (\text{kf-map } g \ \mathfrak{F}) = \text{kf-map } (\text{map-sum } f \ g) (\text{kf-plus } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

2.8 Composition

lemma *kf-comp-dependent-raw-norm-aux*:
fixes $\mathfrak{E} :: \langle 'a \Rightarrow ('e::\text{chilbert-space}, 'f::\text{chilbert-space}, 'g) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('b::\text{chilbert-space}, 'e, 'a) \text{ kraus-family} \rangle$
assumes $B: \langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \text{kf-norm } (\mathfrak{E} \ x) \leq B \rangle$
assumes $[\text{simp}]: \langle B \geq 0 \rangle$
assumes $\langle \text{finite } C \rangle$
assumes $C\text{-subset}: \langle C \subseteq (\lambda((F,y), (E,x)). (E \ o_{CL} \ F, (F,E,y,x))) \text{ ' } (\text{SIGMA } (F,y):\text{Rep-kraus-family } \mathfrak{F}. \text{Rep-kraus-family } (\mathfrak{E} \ y)) \rangle$
shows $\langle (\sum (E,x) \in C. E * \ o_{CL} \ E) \leq (B * \text{kf-norm } \mathfrak{F}) *_{R} \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lift-definition *kf-comp-dependent-raw* :: $\langle ('x \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'y) \text{ kraus-family}) \Rightarrow ('a::\text{chilbert-space}, 'b, 'x) \text{ kraus-family} \rangle$
 $\Rightarrow ('a, 'c, ('a \Rightarrow_{CL} 'b) \times ('b \Rightarrow_{CL} 'c) \times 'x \times 'y) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda \mathfrak{E} \ \mathfrak{F}. \text{if bdd-above } ((\text{kf-norm } \circ \ \mathfrak{E}) \text{ ' } \text{kf-domain } \mathfrak{F}) \text{ then}$
 $\text{Set.filter } (\lambda(EF, -). EF \neq 0) ((\lambda((F,y), (E::'b \Rightarrow_{CL} 'c, x::'y)). (E \ o_{CL} \ F, (F,E,y,x))) \text{ ' } (\text{SIGMA}$
 $(F::'a \Rightarrow_{CL} 'b, y::'x):\text{Rep-kraus-family } \mathfrak{F}. (\text{Rep-kraus-family } (\mathfrak{E} \ y)))$
 $\text{else } \{\} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-norm-leq*:
fixes $\mathfrak{E} :: \langle 'a \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'd) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('e::\text{chilbert-space}, 'b, 'a) \text{ kraus-family} \rangle$
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \text{kf-norm } (\mathfrak{E} \ x) \leq B \rangle$
assumes $\langle B \geq 0 \rangle$
shows $\langle \text{kf-norm } (\text{kf-comp-dependent-raw } \mathfrak{E} \ \mathfrak{F}) \leq B * \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

hide-fact *kf-comp-dependent-raw-norm-aux*

definition $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} = \text{kf-map } (\lambda(F,E,y,x). (y,x)) (\text{kf-comp-dependent-raw } \mathfrak{E} \ \mathfrak{F}) \rangle$

definition $\langle \text{kf-comp } \mathfrak{E} \ \mathfrak{F} = \text{kf-comp-dependent } (\lambda\cdot. \mathfrak{E}) \ \mathfrak{F} \rangle$

lemma *kf-comp-dependent-norm-leq*:

assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \text{kf-norm } (\mathfrak{E} \ x) \leq B \rangle$
assumes $\langle B \geq 0 \rangle$
shows $\langle \text{kf-norm } (\text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F}) \leq B * \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-norm-leq*:

shows $\langle \text{kf-norm } (\text{kf-comp } \mathfrak{E} \ \mathfrak{F}) \leq \text{kf-norm } \mathfrak{E} * \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-apply*:

fixes $\mathfrak{E} :: \langle 'y \Rightarrow ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('c::\text{chilbert-space}, 'a, 'y) \text{ kraus-family} \rangle$
assumes $\langle \text{bdd-above } ((\text{kf-norm } \circ \mathfrak{E}) \ \text{kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent-raw } \mathfrak{E} \ \mathfrak{F} *_{kr} \varrho$
 $= (\sum_{\infty} (F, y) \in \text{Rep-kraus-family } \mathfrak{F}. \ \mathfrak{E} \ y *_{kr} \text{ sandwich-tc } F \ \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-apply*:

fixes $\mathfrak{E} :: \langle 'y \Rightarrow ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('c::\text{chilbert-space}, 'a, 'y) \text{ kraus-family} \rangle$
assumes $\langle \text{bdd-above } ((\text{kf-norm } \circ \mathfrak{E}) \ \text{kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} *_{kr} \varrho$
 $= (\sum_{\infty} (F, y) \in \text{Rep-kraus-family } \mathfrak{F}. \ \mathfrak{E} \ y *_{kr} \text{ sandwich-tc } F \ \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-apply*:

shows $\langle \text{kf-apply } (\text{kf-comp } \mathfrak{E} \ \mathfrak{F}) = \text{kf-apply } \mathfrak{E} \circ \text{kf-apply } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-cong-weak*: $\langle \text{kf-comp } F \ G =_{kr} \text{kf-comp } F' \ G' \rangle$

if $\langle F =_{kr} F' \rangle$ **and** $\langle G =_{kr} G' \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-assoc*:

fixes $\mathfrak{E} :: \langle 'f \Rightarrow ('c::\text{chilbert-space}, 'd::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle 'g \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'g) \text{ kraus-family} \rangle$
defines $\langle \text{reorder} :: 'a \Rightarrow_{CL} 'c \times 'c \Rightarrow_{CL} 'd \times ('a \Rightarrow_{CL} 'b \times 'b \Rightarrow_{CL} 'c \times 'g \times 'f) \times 'e$
 $\Rightarrow 'a \Rightarrow_{CL} 'b \times 'b \Rightarrow_{CL} 'd \times 'g \times 'b \Rightarrow_{CL} 'c \times 'c \Rightarrow_{CL} 'd \times 'f \times 'e \equiv$
 $\lambda(FG::'a \Rightarrow_{CL} 'c, E::'c \Rightarrow_{CL} 'd, (G::'a \Rightarrow_{CL} 'b, F::'b \Rightarrow_{CL} 'c, g::'g, f::'f), e::'e).$
 $(G, E \ o_{CL} \ F, g, F, E, f, e) \rangle$
assumes $\langle \text{bdd-above } (\text{range } (\text{kf-norm } \circ \mathfrak{E})) \rangle$
assumes $\langle \text{bdd-above } (\text{range } (\text{kf-norm } \circ \mathfrak{F})) \rangle$
shows $\langle \text{kf-comp-dependent-raw } (\lambda g::'g. \text{kf-comp-dependent-raw } \mathfrak{E} \ (\mathfrak{F} \ g)) \ \mathfrak{G}$
 $= \text{kf-map-inj reorder } (\text{kf-comp-dependent-raw } (\lambda(-, -, -, f). \ \mathfrak{E} \ f) \ (\text{kf-comp-dependent-raw } \mathfrak{F}$

$\mathfrak{G})\rangle$
 (is $\langle ?lhs = ?rhs \rangle$)
 $\langle proof \rangle$

lemma *kf-filter-comp-dependent*:

fixes $\mathfrak{F} :: \langle 'e \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{F}) \text{ ' kf-domain } \mathfrak{E}) \rangle$
shows $\langle \text{kf-filter } (\lambda(e,f). F \ e \ f \wedge E \ e) \ (\text{kf-comp-dependent } \mathfrak{F} \ \mathfrak{E})$
 $= \text{kf-comp-dependent } (\lambda e. \text{kf-filter } (F \ e) \ (\mathfrak{F} \ e)) \ (\text{kf-filter } E \ \mathfrak{E}) \rangle$
 $\langle proof \rangle$

lemma *kf-comp-assoc-weak*:

fixes $\mathfrak{E} :: \langle ('c::\text{chilbert-space}, 'd::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'g) \text{ kraus-family} \rangle$
shows $\langle \text{kf-comp } (\text{kf-comp } \mathfrak{E} \ \mathfrak{F}) \ \mathfrak{G} =_{kr} \text{kf-comp } \mathfrak{E} \ (\text{kf-comp } \mathfrak{F} \ \mathfrak{G}) \rangle$
 $\langle proof \rangle$

lemma *kf-comp-dependent-raw-cong-left*:

assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ ' kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}') \text{ ' kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \bigwedge x. x \in \text{snd } \text{' Rep-kraus-family } \mathfrak{F} \implies \mathfrak{E} \ x = \mathfrak{E}' \ x \rangle$
shows $\langle \text{kf-comp-dependent-raw } \mathfrak{E} \ \mathfrak{F} = \text{kf-comp-dependent-raw } \mathfrak{E}' \ \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-comp-dependent-cong-left*:

assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ ' kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}') \text{ ' kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \mathfrak{E} \ x = \mathfrak{E}' \ x \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} = \text{kf-comp-dependent } \mathfrak{E}' \ \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *kf-domain-comp-dependent-raw-subset*:

$\langle \text{kf-domain } (\text{kf-comp-dependent-raw } \mathfrak{E} \ \mathfrak{F}) \subseteq \text{UNIV} \times \text{UNIV} \times (\text{SIGMA } x:\text{kf-domain } \mathfrak{F}. \text{kf-domain } (\mathfrak{E} \ x)) \rangle$
 $\langle proof \rangle$

lemma *kf-domain-comp-dependent-subset*:

$\langle \text{kf-domain } (\text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F}) \subseteq (\text{SIGMA } x:\text{kf-domain } \mathfrak{F}. \text{kf-domain } (\mathfrak{E} \ x)) \rangle$
 $\langle proof \rangle$

lemma *kf-domain-comp-subset*: $\langle \text{kf-domain } (\text{kf-comp } \mathfrak{E} \ \mathfrak{F}) \subseteq \text{kf-domain } \mathfrak{F} \times \text{kf-domain } \mathfrak{E} \rangle$

$\langle proof \rangle$

lemma *kf-apply-comp-dependent-cong*:

fixes $\mathfrak{E} :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'e1) \text{ kraus-family} \rangle$

and $\mathfrak{E}' :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'e2) \text{ kraus-family} \rangle$
and $\mathfrak{F} \mathfrak{F}' :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
assumes $\text{bdd} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ ' kf-domain } \mathfrak{F}) \rangle$
assumes $\text{bdd}' :: \langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}') \text{ ' kf-domain } \mathfrak{F}') \rangle$
assumes $\langle f \in \text{kf-domain } \mathfrak{F} \implies \text{kf-apply-on } (\mathfrak{E} \ f) \ E = \text{kf-apply-on } (\mathfrak{E}' \ f) \ E' \rangle$
assumes $\langle \text{kf-apply-on } \mathfrak{F} \ \{f\} = \text{kf-apply-on } \mathfrak{F}' \ \{f\} \rangle$
shows $\langle \text{kf-apply-on } (\text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F}) \ (\{f\} \times E) = \text{kf-apply-on } (\text{kf-comp-dependent } \mathfrak{E}' \ \mathfrak{F}') \ (\{f\} \times E') \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-cong-weak*:

fixes $\mathfrak{E} :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'e1) \text{ kraus-family} \rangle$
and $\mathfrak{E}' :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'e2) \text{ kraus-family} \rangle$
and $\mathfrak{F} \mathfrak{F}' :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
assumes $\text{bdd} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ ' kf-domain } \mathfrak{F}) \rangle$
assumes $\text{eq} :: \langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \mathfrak{E} \ x =_{kr} \mathfrak{E}' \ x \rangle$
assumes $\langle \mathfrak{F} \equiv_{kr} \mathfrak{F}' \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} =_{kr} \text{kf-comp-dependent } \mathfrak{E}' \ \mathfrak{F}' \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-assoc*:

fixes $\mathfrak{E} :: \langle 'g \Rightarrow 'f \Rightarrow ('c::\text{chilbert-space}, 'd::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle 'g \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'g) \text{ kraus-family} \rangle$
assumes $\text{bdd-E} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \text{case-prod } \mathfrak{E}) \text{ ' } (\text{SIGMA } x:\text{kf-domain } \mathfrak{G}. \text{kf-domain } (\mathfrak{F} \ x))) \rangle$
assumes $\text{bdd-F} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{F}) \text{ ' kf-domain } \mathfrak{G}) \rangle$
shows $\langle (\text{kf-comp-dependent } (\lambda g. \text{kf-comp-dependent } (\mathfrak{E} \ g) (\mathfrak{F} \ g)) \ \mathfrak{G}) \equiv_{kr} \text{kf-map } (\lambda((g,f),e). (g,f,e)) (\text{kf-comp-dependent } (\lambda(g,f). \mathfrak{E} \ g \ f) (\text{kf-comp-dependent } \mathfrak{F} \ \mathfrak{G})) \rangle$
 $\langle \text{is } \langle ?lhs \equiv_{kr} ?rhs \rangle \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-assoc-weak*:

fixes $\mathfrak{E} :: \langle 'g \Rightarrow 'f \Rightarrow ('c::\text{chilbert-space}, 'd::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle 'g \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'g) \text{ kraus-family} \rangle$
assumes $\text{bdd-E} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \text{case-prod } \mathfrak{E}) \text{ ' } (\text{SIGMA } x:\text{kf-domain } \mathfrak{G}. \text{kf-domain } (\mathfrak{F} \ x))) \rangle$
assumes $\text{bdd-F} :: \langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{F}) \text{ ' kf-domain } \mathfrak{G}) \rangle$
shows $\langle \text{kf-comp-dependent } (\lambda g. \text{kf-comp-dependent } (\mathfrak{E} \ g) (\mathfrak{F} \ g)) \ \mathfrak{G} =_{kr} \text{kf-comp-dependent } (\lambda(g,f). \mathfrak{E} \ g \ f) (\text{kf-comp-dependent } \mathfrak{F} \ \mathfrak{G}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-comp-assoc-weak*:

fixes $\mathfrak{E} :: \langle ('c::\text{chilbert-space}, 'd::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle 'g \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$

and $\mathfrak{G} :: \langle ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'g) \text{ kraus-family} \rangle$
assumes $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{F}) \text{ 'kf-domain } \mathfrak{G}) \rangle$
shows $\langle \text{kf-comp-dependent } (\lambda g. \text{kf-comp } \mathfrak{E} (\mathfrak{F} \ g)) \ \mathfrak{G} =_{kr}$
 $\text{kf-comp } \mathfrak{E} (\text{kf-comp-dependent } \mathfrak{F} \ \mathfrak{G}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-comp-dependent-assoc-weak:*

fixes $\mathfrak{E} :: \langle ('f \Rightarrow ('c::\text{hilbert-space}, 'd::\text{hilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('b::\text{hilbert-space}, 'c::\text{hilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'g) \text{ kraus-family} \rangle$
assumes *bdd-E*: $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp } (\text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F}) \ \mathfrak{G} =_{kr}$
 $\text{kf-comp-dependent } (\lambda(-,f). \ \mathfrak{E} \ f) (\text{kf-comp } \mathfrak{F} \ \mathfrak{G}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-assoc:*

fixes $\mathfrak{E} :: \langle ('c::\text{hilbert-space}, 'd::\text{hilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} :: \langle ('b::\text{hilbert-space}, 'c::\text{hilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{G} :: \langle ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'g) \text{ kraus-family} \rangle$
shows $\langle \text{kf-comp } (\text{kf-comp } \mathfrak{E} \ \mathfrak{F}) \ \mathfrak{G} \equiv_{kr}$
 $\text{kf-map } (\lambda((g,f),e). \ (g,f,e)) (\text{kf-comp } \mathfrak{E} (\text{kf-comp } \mathfrak{F} \ \mathfrak{G})) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-cong:*

fixes $\mathfrak{E} \ \mathfrak{E}' :: \langle ('f \Rightarrow ('b::\text{hilbert-space}, 'c::\text{hilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} \ \mathfrak{F}' :: \langle ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'f) \text{ kraus-family} \rangle$
assumes *bdd*: $\langle \text{bdd-above } ((\text{kf-norm } o \ \mathfrak{E}) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \bigwedge x. x \in \text{kf-domain } \mathfrak{F} \implies \mathfrak{E} \ x \equiv_{kr} \mathfrak{E}' \ x \rangle$
assumes $\langle \mathfrak{F} \equiv_{kr} \mathfrak{F}' \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \ \mathfrak{F} \equiv_{kr} \text{kf-comp-dependent } \mathfrak{E}' \ \mathfrak{F}' \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-cong:*

fixes $\mathfrak{E} \ \mathfrak{E}' :: \langle ('b::\text{hilbert-space}, 'c::\text{hilbert-space}, 'e) \text{ kraus-family} \rangle$
and $\mathfrak{F} \ \mathfrak{F}' :: \langle ('a::\text{hilbert-space}, 'b::\text{hilbert-space}, 'f) \text{ kraus-family} \rangle$
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{E}' \rangle$
assumes $\langle \mathfrak{F} \equiv_{kr} \mathfrak{F}' \rangle$
shows $\langle \text{kf-comp } \mathfrak{E} \ \mathfrak{F} \equiv_{kr} \text{kf-comp } \mathfrak{E}' \ \mathfrak{F}' \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-comp-dependent-raw-of-op:*

shows $\langle \text{kf-bound } (\text{kf-comp-dependent-raw } \mathfrak{E} (\text{kf-of-op } U))$
 $= \text{sandwich } (U^*) (\text{kf-bound } (\mathfrak{E} ())) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-comp-dependent-of-op:*

shows $\langle \text{kf-bound } (\text{kf-comp-dependent } \mathfrak{E} (\text{kf-of-op } U)) = \text{sandwich } (U^*) (\text{kf-bound } (\mathfrak{E} ())) \rangle$

$\langle proof \rangle$

lemma *kf-bound-comp-of-op*:

shows $\langle kf-bound (kf-comp \mathfrak{E} (kf-of-op U)) = sandwich (U^*) (kf-bound \mathfrak{E}) \rangle$

$\langle proof \rangle$

lemma *kf-norm-comp-dependent-of-op-coiso*:

assumes $\langle isometry (U^*) \rangle$

shows $\langle kf-norm (kf-comp-dependent \mathfrak{E} (kf-of-op U)) = kf-norm (\mathfrak{E} ()) \rangle$

$\langle proof \rangle$

lemma *kf-norm-comp-of-op-coiso*:

assumes $\langle isometry (U^*) \rangle$

shows $\langle kf-norm (kf-comp \mathfrak{E} (kf-of-op U)) = kf-norm \mathfrak{E} \rangle$

$\langle proof \rangle$

lemma *kf-bound-comp-dependend-raw-iso*:

assumes $\langle isometry U \rangle$

shows $\langle kf-bound (kf-comp-dependent-raw (\lambda-. kf-of-op U) \mathfrak{E}) = kf-bound \mathfrak{E} \rangle$

$\langle proof \rangle$

lemma *kf-bound-comp-dependent-iso*:

assumes $\langle isometry U \rangle$

shows $\langle kf-bound (kf-comp-dependent (\lambda-. kf-of-op U) \mathfrak{E}) = kf-bound \mathfrak{E} \rangle$

$\langle proof \rangle$

lemma *kf-bound-comp-iso*:

assumes $\langle isometry U \rangle$

shows $\langle kf-bound (kf-comp (kf-of-op U) \mathfrak{E}) = kf-bound \mathfrak{E} \rangle$

$\langle proof \rangle$

lemma *kf-norm-comp-dependent-iso*:

assumes $\langle isometry U \rangle$

shows $\langle kf-norm (kf-comp-dependent (\lambda-. kf-of-op U) \mathfrak{E}) = kf-norm \mathfrak{E} \rangle$

$\langle proof \rangle$

lemma *kf-norm-comp-iso*:

assumes $\langle isometry U \rangle$

shows $\langle kf-norm (kf-comp (kf-of-op U) \mathfrak{E}) = kf-norm \mathfrak{E} \rangle$

$\langle proof \rangle$

lemma *kf-comp-dependent-raw-0-right[simp]*: $\langle kf-comp-dependent-raw \mathfrak{E} 0 = 0 \rangle$

$\langle proof \rangle$

lemma *kf-comp-dependent-raw-0-left[simp]*: $\langle kf-comp-dependent-raw 0 \mathfrak{E} = 0 \rangle$

$\langle proof \rangle$

lemma *kf-comp-dependent-0-left[simp]*: $\langle \text{kf-comp-dependent } (\lambda-. 0) E = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-0-right[simp]*: $\langle \text{kf-comp-dependent } E 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-0-left[simp]*: $\langle \text{kf-comp } 0 E = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-0-right[simp]*: $\langle \text{kf-comp } E 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-comp*:

fixes $\mathfrak{F} :: \langle ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'f) \text{ kraus-family} \rangle$
and $\mathfrak{E} :: \langle ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$
shows $\langle \text{kf-filter } (\lambda(e,f). F f \wedge E e) (\text{kf-comp } \mathfrak{F} \mathfrak{E})$
 $= \text{kf-comp } (\text{kf-filter } F \mathfrak{F}) (\text{kf-filter } E \mathfrak{E}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-invalid*:

assumes $\langle \neg \text{bdd-above } ((\text{kf-norm } \circ \mathfrak{E}) \text{ `kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \mathfrak{F} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-map-left*:

$\langle \text{kf-comp-dependent } (\lambda x. \text{kf-map } (f x) (E x)) F$
 $\equiv_{kr} \text{kf-map } (\lambda(x,y). (x, f x y)) (\text{kf-comp-dependent } E F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-map-right*:

$\langle \text{kf-comp-dependent } E (\text{kf-map } f F)$
 $\equiv_{kr} \text{kf-map } (\lambda(x,y). (f x, y)) (\text{kf-comp-dependent } (\lambda x. E (f x)) F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-map-inj-right*:

$\langle \text{kf-comp-dependent-raw } E (\text{kf-map-inj } f F)$
 $= \text{kf-map-inj } (\lambda(E,F,x,y). (E, F, f x, y)) (\text{kf-comp-dependent-raw } (\lambda x. E (f x)) F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-map-inj-right*:

assumes $\langle \text{inj-on } f (\text{kf-domain } F) \rangle$
shows $\langle \text{kf-comp-dependent } E (\text{kf-map-inj } f F)$
 $= \text{kf-map-inj } (\lambda(x,y). (f x, y)) (\text{kf-comp-dependent } (\lambda x. E (f x)) F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-map-right-weak*:

$\langle \text{kf-comp-dependent } E \text{ (kf-map } f \text{ } F) \rangle$
 $=_{k_r} \text{kf-comp-dependent } (\lambda x. E \text{ (} f \text{ } x)) \text{ } F \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-map-left*:
 $\langle \text{kf-comp (kf-map } f \text{ } E) \text{ } F \equiv_{k_r} \text{kf-map } (\lambda(x,y). (x, f \text{ } y)) \text{ (kf-comp } E \text{ } F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-map-right*:
 $\langle \text{kf-comp } E \text{ (kf-map } f \text{ } F) \equiv_{k_r} \text{kf-map } (\lambda(x,y). (f \text{ } x, y)) \text{ (kf-comp } E \text{ } F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-map-both*:
 $\langle \text{kf-comp (kf-map } e \text{ } E) \text{ (kf-map } f \text{ } F) \equiv_{k_r} \text{kf-map } (\lambda(x,y). (f \text{ } x, e \text{ } y)) \text{ (kf-comp } E \text{ } F) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-apply-commute*:
assumes $\langle \text{kf-operators } \mathfrak{F} \subseteq \text{commutant (kf-operators } \mathfrak{E}) \rangle$
shows $\langle \text{kf-apply } \mathfrak{F} \circ \text{kf-apply } \mathfrak{E} = \text{kf-apply } \mathfrak{E} \circ \text{kf-apply } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-commute-weak*:
assumes $\langle \text{kf-operators } \mathfrak{F} \subseteq \text{commutant (kf-operators } \mathfrak{E}) \rangle$
shows $\langle \text{kf-comp } \mathfrak{F} \text{ } \mathfrak{E} =_{k_r} \text{kf-comp } \mathfrak{E} \text{ } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-commute*:
assumes $\langle \text{kf-operators } \mathfrak{F} \subseteq \text{commutant (kf-operators } \mathfrak{E}) \rangle$
shows $\langle \text{kf-comp } \mathfrak{F} \text{ } \mathfrak{E} \equiv_{k_r} \text{kf-map prod.swap (kf-comp } \mathfrak{E} \text{ } \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-apply-on-singleton*:
 $\langle \text{kf-comp } \mathfrak{E} \text{ } \mathfrak{F} *_{k_r} @\{x\} \varrho = \mathfrak{E} *_{k_r} @\{\text{snd } x\} (\mathfrak{F} *_{k_r} @\{\text{fst } x\} \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-apply-on-singleton*:
assumes $\langle \text{bdd-above } ((\text{kf-norm } \circ \mathfrak{E}) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent } \mathfrak{E} \text{ } \mathfrak{F} *_{k_r} @\{x\} \varrho = \mathfrak{E} (\text{fst } x) *_{k_r} @\{\text{snd } x\} (\mathfrak{F} *_{k_r} @\{\text{fst } x\} \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-id-left*: $\langle \text{kf-comp kf-id } \mathfrak{E} \equiv_{k_r} \text{kf-map } (\lambda x. (x, ())) \text{ } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-id-right*: $\langle \text{kf-comp } \mathfrak{E} \text{ kf-id} \equiv_{k_r} \text{kf-map } (\lambda x. ((), x)) \text{ } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-raw-kf-plus-left*:
fixes $\mathfrak{D} :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'd) \text{ kraus-family} \rangle$
fixes $\mathfrak{E} :: \langle 'f \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'e) \text{ kraus-family} \rangle$

fixes $\mathfrak{F} :: \langle ('a::\text{chilbert-space}, 'b, 'f) \text{ kraus-family} \rangle$
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{D} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{E} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent-raw } (\lambda x. \text{kf-plus } (\mathfrak{D} x) (\mathfrak{E} x)) \mathfrak{F} =$
 $\text{kf-map-inj } (\lambda x. \text{case } x \text{ of Inl } (F, D, f, d) \Rightarrow (F, D, f, \text{Inl } d) \mid \text{Inr } (F, E, f, e) \Rightarrow (F, E, f, \text{Inr } e))$
 $(\text{kf-plus } (\text{kf-comp-dependent-raw } \mathfrak{D} \mathfrak{F}) (\text{kf-comp-dependent-raw } \mathfrak{E} \mathfrak{F})) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-kf-plus-left:*

assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{D} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{E} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent } (\lambda x. \text{kf-plus } (\mathfrak{D} x) (\mathfrak{E} x)) \mathfrak{F} =$
 $\text{kf-map-inj } (\lambda x. \text{case } x \text{ of Inl } (f, d) \Rightarrow (f, \text{Inl } d) \mid \text{Inr } (f, e) \Rightarrow (f, \text{Inr } e)) (\text{kf-plus } (\text{kf-comp-dependent } \mathfrak{D} \mathfrak{F}) (\text{kf-comp-dependent } \mathfrak{E} \mathfrak{F})) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-map-inj-twice:*

shows $\langle \text{kf-map-inj } f (\text{kf-map-inj } g \mathfrak{E}) = \text{kf-map-inj } (f \circ g) \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-kf-plus-left':*

assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{D} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{E} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-plus } (\text{kf-comp-dependent } \mathfrak{D} \mathfrak{F}) (\text{kf-comp-dependent } \mathfrak{E} \mathfrak{F}) =$
 $\text{kf-map-inj } (\lambda (f, de). \text{case } de \text{ of Inl } d \Rightarrow \text{Inl } (f, d) \mid \text{Inr } e \Rightarrow \text{Inr } (f, e)) (\text{kf-comp-dependent } (\lambda x. \text{kf-plus } (\mathfrak{D} x) (\mathfrak{E} x)) \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-comp-dependent-plus-left:*

assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{D} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
assumes $\langle \text{bdd-above } ((\lambda x. \text{kf-norm } (\mathfrak{E} x)) \text{ 'kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-comp-dependent } (\lambda x. \mathfrak{D} x + \mathfrak{E} x) \mathfrak{F} \equiv_{kr} \text{kf-comp-dependent } \mathfrak{D} \mathfrak{F} + \text{kf-comp-dependent } \mathfrak{E} \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

2.9 Infinite sums

lift-definition *kf-infsum* :: $\langle ('a \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'x) \text{ kraus-family}) \Rightarrow 'a \text{ set} \Rightarrow ('b, 'c, 'a \times 'x) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda \mathfrak{E} A. \text{if summable-on-in cweak-operator-topology } (\lambda a. \text{kf-bound } (\mathfrak{E} a)) A$
 $\text{then } (\lambda (a, (E, x)). (E, (a, x))) \text{ 'Sigma } A (\lambda a. \text{Rep-kraus-family } (\mathfrak{E} a)) \text{ else } \{\} \rangle$
 $\langle \text{proof} \rangle$

definition *kf-summable* :: $\langle ('a \Rightarrow ('b::\text{chilbert-space}, 'c::\text{chilbert-space}, 'x) \text{ kraus-family}) \Rightarrow 'a \text{ set} \Rightarrow \text{bool} \rangle$ **where**

$\langle \text{kf-summable } \mathfrak{E} A \longleftrightarrow \text{summable-on-in cweak-operator-topology } (\lambda a. \text{kf-bound } (\mathfrak{E} a)) A \rangle$

lemma *kf-summable-from-abs-summable*:
assumes $\text{sum}: \langle (\lambda a. \text{kf-norm } (\mathfrak{E} a)) \text{ summable-on } A \rangle$
shows $\langle \text{kf-summable } \mathfrak{E} A \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-infsum-apply*:
assumes $\langle \text{kf-summable } \mathfrak{E} A \rangle$
shows $\langle \text{kf-infsum } \mathfrak{E} A *_{kr} \varrho = (\sum_{\infty} a \in A. \mathfrak{E} a *_{kr} \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-infsum-apply-summable*:
assumes $\langle \text{kf-summable } \mathfrak{E} A \rangle$
shows $\langle (\lambda a. \mathfrak{E} a *_{kr} \varrho) \text{ summable-on } A \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-infsum*:
fixes $f :: \langle 'a \Rightarrow ('b::\text{hilbert-space}, 'c::\text{hilbert-space}, 'x) \text{ kraus-family} \rangle$
assumes $\langle \text{kf-summable } f A \rangle$
shows $\langle \text{kf-bound } (\text{kf-infsum } f A) = \text{infsum-in cweak-operator-topology } (\lambda a. \text{kf-bound } (f a)) A \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-infsum*:
assumes $\text{sum}: \langle (\lambda a. \text{kf-norm } (\mathfrak{E} a)) \text{ summable-on } A \rangle$
shows $\langle \text{kf-norm } (\text{kf-infsum } \mathfrak{E} A) \leq (\sum_{\infty} a \in A. \text{kf-norm } (\mathfrak{E} a)) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-infsum*:
assumes $\langle \text{kf-summable } \mathfrak{E} A \rangle$
shows $\langle \text{kf-filter } P (\text{kf-infsum } \mathfrak{E} A) = \text{kf-infsum } (\lambda x. \text{kf-filter } (\lambda x. P (a,x)) (\mathfrak{E} a)) \{a \in A. \exists x. P (a,x)\} \rangle$
 $\langle \text{is } \langle ?lhs = ?rhs \rangle \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-infsum-empty[simp]*: $\langle \text{kf-infsum } \mathfrak{E} \{\} = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-infsum-singleton[simp]*: $\langle \text{kf-infsum } \mathfrak{E} \{a\} = \text{kf-map-inj } (\lambda x. (a,x)) (\mathfrak{E} a) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-infsum-invalid*:
assumes $\langle \neg \text{kf-summable } \mathfrak{E} A \rangle$
shows $\langle \text{kf-infsum } \mathfrak{E} A = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-infsum-cong*:
fixes $\mathfrak{E} \mathfrak{F} :: \langle 'a \Rightarrow ('b::\text{hilbert-space}, 'c::\text{hilbert-space}, 'x) \text{ kraus-family} \rangle$

assumes $\langle \bigwedge a. a \in A \implies \mathfrak{E} a \equiv_{k_r} \mathfrak{F} a \rangle$
shows $\langle kf\text{-infsup } \mathfrak{E} A \equiv_{k_r} kf\text{-infsup } \mathfrak{F} A \rangle$
 $\langle proof \rangle$

2.10 Trace-preserving maps

definition $\langle kf\text{-trace-preserving } \mathfrak{E} \longleftrightarrow (\forall \varrho. trace\text{-tc } (\mathfrak{E} *_{k_r} \varrho) = trace\text{-tc } \varrho) \rangle$

definition $\langle kf\text{-trace-reducing } \mathfrak{E} \longleftrightarrow (\forall \varrho \geq 0. trace\text{-tc } (\mathfrak{E} *_{k_r} \varrho) \leq trace\text{-tc } \varrho) \rangle$

lemma $kf\text{-trace-reducing-iff-norm-leq1}$: $\langle kf\text{-trace-reducing } \mathfrak{E} \longleftrightarrow kf\text{-norm } \mathfrak{E} \leq 1 \rangle$
 $\langle proof \rangle$

lemma $kf\text{-trace-preserving-iff-bound-id}$: $\langle kf\text{-trace-preserving } \mathfrak{E} \longleftrightarrow kf\text{-bound } \mathfrak{E} = id\text{-cblinfun} \rangle$
 $\langle proof \rangle$

lemma $kf\text{-trace-norm-preserving}$: $\langle kf\text{-norm } \mathfrak{E} \leq 1 \rangle$ **if** $\langle kf\text{-trace-preserving } \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma $kf\text{-trace-norm-preserving-eq}$:
fixes $\mathfrak{E} :: \langle ('a :: \{chilbert\text{-space}, not\text{-singleton}\}, 'b :: chilbert\text{-space}, 'c) \text{ kraus-family} \rangle$
assumes $\langle kf\text{-trace-preserving } \mathfrak{E} \rangle$
shows $\langle kf\text{-norm } \mathfrak{E} = 1 \rangle$
 $\langle proof \rangle$

lemma $kf\text{-trace-preserving-map[simp]}$: $\langle kf\text{-trace-preserving } (kf\text{-map } f \ \mathfrak{E}) \longleftrightarrow kf\text{-trace-preserving } \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma $kf\text{-trace-reducing-map[simp]}$: $\langle kf\text{-trace-reducing } (kf\text{-map } f \ \mathfrak{E}) \longleftrightarrow kf\text{-trace-reducing } \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma $kf\text{-trace-preserving-id[iff]}$: $\langle kf\text{-trace-preserving } kf\text{-id} \rangle$
 $\langle proof \rangle$

lemma $kf\text{-trace-reducing-id[iff]}$: $\langle kf\text{-trace-reducing } kf\text{-id} \rangle$
 $\langle proof \rangle$

2.11 Sampling

lift-definition $kf\text{-sample}$:: $\langle ('x \implies real) \implies ('a :: chilbert\text{-space}, 'a, 'x) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda p. \text{if } (\forall x. p \ x \geq 0) \wedge p \text{ summable-on UNIV} \text{ then } Set.filter (\lambda(E, -). E \neq 0) (\text{range } (\lambda x. (\text{sqrt } (p \ x) *_{k_r} id\text{-cblinfun}, x))) \text{ else } \{\} \rangle$
 $\langle proof \rangle$

lemma $kf\text{-sample-norm}$:
fixes $p :: \langle 'x \implies real \rangle$
assumes $\langle \bigwedge x. p \ x \geq 0 \rangle$
assumes $\langle p \text{ summable-on UNIV} \rangle$

shows $\langle \text{kf-norm } (\text{kf-sample } p :: ('a::\{\text{chilbert-space, not-singleton}\}, 'a, 'x) \text{ kraus-family})$
 $= (\sum_{\infty} x. p \ x) \rangle$
 $\langle \text{proof} \rangle$

2.12 Trace

lift-definition $\text{kf-trace} :: \langle 'a \text{ set} \Rightarrow ('a::\text{chilbert-space}, 'b::\text{one-dim}, 'a) \text{ kraus-family} \rangle$ **is**
 $\langle \lambda B. \text{ if is-onb } B \text{ then } (\lambda x. (\text{vector-to-cblinfun } x^*, x)) \text{ ' } B \text{ else } \{\} \rangle$
 $\langle \text{proof} \rangle$

lemma kf-trace-is-trace :
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-trace } B *_{kr} \varrho = \text{one-dim-iso } (\text{trace-tc } \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{kf-eq-weak-kf-trace}$:
assumes $\langle \text{is-onb } A \rangle$ **and** $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-trace } A =_{kr} \text{kf-trace } B \rangle$
 $\langle \text{proof} \rangle$

lemma trace-is-kf-trace :
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{trace-tc } t = \text{one-dim-iso } (\text{kf-trace } B *_{kr} t) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{kf-trace-bound[simp]}$:
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-bound } (\text{kf-trace } B) = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{kf-trace-norm-eq1[simp]}$:
fixes $B :: \langle 'a::\{\text{chilbert-space}, \text{not-singleton}\} \text{ set} \rangle$
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-norm } (\text{kf-trace } B) = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{kf-trace-norm-leq1[simp]}$:
fixes $B :: \langle 'a::\text{chilbert-space} \text{ set} \rangle$
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-norm } (\text{kf-trace } B) \leq 1 \rangle$
 $\langle \text{proof} \rangle$

2.13 Constant maps

lift-definition $\text{kf-constant-onedim} :: \langle ('b, 'b) \text{ trace-class} \Rightarrow ('a::\text{one-dim}, 'b::\text{chilbert-space}, \text{unit})$
 $\text{kraus-family} \rangle$ **is**
 $\langle \lambda t::('b, 'b) \text{ trace-class. if } t \geq 0 \text{ then}$
 $\text{Set.filter } (\lambda(E, -). E \neq 0) ((\lambda v. (\text{vector-to-cblinfun } v, ())) \text{ ' spectral-dec-vecs-tc } t)$
 $\text{else } \{\} \rangle$

⟨proof⟩

definition *kf-constant* :: ⟨('b,'b) trace-class ⇒ ('a::chilbert-space, 'b::chilbert-space, unit) kraus-family⟩
where

⟨*kf-constant* ρ = *kf-flatten* (*kf-comp* (*kf-constant-onedim* ρ :: (complex,-,-) kraus-family) (*kf-trace* some-chilbert-basis))⟩

lemma *kf-constant-onedim-invalid*: ⟨¬ ρ ≥ 0 ⇒ *kf-constant-onedim* ρ = 0⟩
⟨proof⟩

lemma *kf-constant-invalid*: ⟨¬ ρ ≥ 0 ⇒ *kf-constant* ρ = 0⟩
⟨proof⟩

lemma *kf-constant-onedim-apply*:
assumes ⟨ρ ≥ 0⟩
shows ⟨*kf-apply* (*kf-constant-onedim* ρ) σ = *one-dim-iso* σ *_C ρ⟩
⟨proof⟩

lemma *kf-constant-apply*:
assumes ⟨ρ ≥ 0⟩
shows ⟨*kf-apply* (*kf-constant* ρ) σ = *trace-tc* σ *_C ρ⟩
⟨proof⟩

lemma *kf-bound-constant-onedim[simp]*:
fixes ρ :: ⟨('a::chilbert-space, 'a) trace-class⟩
assumes ⟨ρ ≥ 0⟩
shows ⟨*kf-bound* (*kf-constant-onedim* ρ) = *norm* ρ *_R *id-cblinfun*⟩
⟨proof⟩

lemma *kf-bound-constant[simp]*:
fixes ρ :: ⟨('a::chilbert-space, 'a) trace-class⟩
assumes ⟨ρ ≥ 0⟩
shows ⟨*kf-bound* (*kf-constant* ρ) = *norm* ρ *_R *id-cblinfun*⟩
⟨proof⟩

lemma *kf-norm-constant-onedim[simp]*:
assumes ⟨ρ ≥ 0⟩
shows ⟨*kf-norm* (*kf-constant-onedim* ρ) = *norm* ρ⟩
⟨proof⟩

lemma *kf-norm-constant*:
assumes ⟨ρ ≥ 0⟩
shows ⟨*kf-norm* (*kf-constant* ρ :: ('a::{chilbert-space,not-singleton}, 'b::chilbert-space,-) kraus-family) = *norm* ρ⟩
⟨proof⟩

lemma *kf-norm-constant-leq*:
shows ⟨*kf-norm* (*kf-constant* ρ) ≤ *norm* ρ⟩
⟨proof⟩

lemma *kf-comp-constant-right*:

assumes $[iff]: \langle t \geq 0 \rangle$

shows $\langle kf\text{-map } fst (kf\text{-comp } E (kf\text{-constant } t)) \equiv_{kr} kf\text{-constant } (E *_{kr} t) \rangle$

$\langle proof \rangle$

lemma *kf-comp-constant-right-weak*:

assumes $[iff]: \langle t \geq 0 \rangle$

shows $\langle kf\text{-comp } E (kf\text{-constant } t) =_{kr} kf\text{-constant } (E *_{kr} t) \rangle$

$\langle proof \rangle$

2.14 Tensor products

lemma *kf-tensor-raw-bound-aux*:

fixes $\mathfrak{E} :: \langle ('a \text{ ell2} \Rightarrow_{CL} 'b \text{ ell2} \times 'x) \text{ set} \rangle$ **and** $\mathfrak{F} :: \langle ('c \text{ ell2} \Rightarrow_{CL} 'd \text{ ell2} \times 'y) \text{ set} \rangle$

assumes $\langle \bigwedge S. \text{finite } S \implies S \subseteq \mathfrak{E} \implies (\sum_{(E,x) \in S} E *_{o_{CL}} E) \leq B \rangle$

assumes $\langle \bigwedge S. \text{finite } S \implies S \subseteq \mathfrak{F} \implies (\sum_{(E,x) \in S} E *_{o_{CL}} E) \leq C \rangle$

assumes $\langle \text{finite } U \rangle$

assumes $\langle U \subseteq ((\lambda((E,x), (F,y)). (E \otimes_o F, E, F, x, y))) '(\mathfrak{E} \times \mathfrak{F}) \rangle$

shows $\langle (\sum_{(G,z) \in U} G *_{o_{CL}} G) \leq B \otimes_o C \rangle$

$\langle proof \rangle$

lift-definition *kf-tensor-raw* :: $\langle ('a \text{ ell2}, 'b \text{ ell2}, 'x) \text{ kraus-family} \Rightarrow ('c \text{ ell2}, 'd \text{ ell2}, 'y) \text{ kraus-family} \Rightarrow$

$\langle ('a \times 'c) \text{ ell2}, ('b \times 'd) \text{ ell2}, ((('a \text{ ell2} \Rightarrow_{CL} 'b \text{ ell2}) \times ('c \text{ ell2} \Rightarrow_{CL} 'd \text{ ell2}) \times 'x \times 'y)) \text{ kraus-family} \rangle$

is

$\langle \lambda \mathfrak{E} \mathfrak{F}. (\lambda((E,x), (F,y)). (E \otimes_o F, (E,F,x,y))) '(\mathfrak{E} \times \mathfrak{F}) \rangle$

$\langle proof \rangle$

lemma *kf-apply-tensor-raw-as-infsum*:

$\langle kf\text{-tensor-raw } \mathfrak{E} \mathfrak{F} *_{kr} \varrho = (\sum_{\infty} ((E,x),(F,y)) \in \text{Rep-kraus-family } \mathfrak{E} \times \text{Rep-kraus-family } \mathfrak{F}. \text{sandwich-tc } (E \otimes_o F) \varrho) \rangle$

$\langle proof \rangle$

lemma *kf-apply-tensor-raw*:

shows $\langle kf\text{-tensor-raw } \mathfrak{E} \mathfrak{F} *_{kr} \text{tc-tensor } \varrho \sigma = \text{tc-tensor } (\mathfrak{E} *_{kr} \varrho) (\mathfrak{F} *_{kr} \sigma) \rangle$

$\langle proof \rangle$

hide-fact *kf-tensor-raw-bound-aux*

definition $\langle kf\text{-tensor } \mathfrak{E} \mathfrak{F} = kf\text{-map } (\lambda(E, F, x, y). (x,y)) (kf\text{-tensor-raw } \mathfrak{E} \mathfrak{F}) \rangle$

lemma *kf-apply-tensor*:

$\langle kf\text{-tensor } \mathfrak{E} \mathfrak{F} *_{kr} \text{tc-tensor } \varrho \sigma = \text{tc-tensor } (\mathfrak{E} *_{kr} \varrho) (\mathfrak{F} *_{kr} \sigma) \rangle$

$\langle proof \rangle$

lemma *kf-apply-tensor-as-infsum*:

$\langle kf\text{-tensor } \mathfrak{E} \mathfrak{F} *_{kr} \varrho = (\sum_{\infty} ((E,x),(F,y)) \in \text{Rep-kraus-family } \mathfrak{E} \times \text{Rep-kraus-family } \mathfrak{F}. \text{sand-}$

wich-tc $\langle E \otimes_o F \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-tensor-raw*:

$\langle \text{kf-bound } (\text{kf-tensor-raw } \mathfrak{E} \mathfrak{F}) = \text{kf-bound } \mathfrak{E} \otimes_o \text{kf-bound } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-tensor*:

$\langle \text{kf-bound } (\text{kf-tensor } \mathfrak{E} \mathfrak{F}) = \text{kf-bound } \mathfrak{E} \otimes_o \text{kf-bound } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-tensor*:

$\langle \text{kf-norm } (\text{kf-tensor } \mathfrak{E} \mathfrak{F}) = \text{kf-norm } \mathfrak{E} * \text{kf-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-cong-weak*:

assumes $\langle \mathfrak{E} =_{kr} \mathfrak{E}' \rangle$
assumes $\langle \mathfrak{F} =_{kr} \mathfrak{F}' \rangle$
shows $\langle \text{kf-tensor } \mathfrak{E} \mathfrak{F} =_{kr} \text{kf-tensor } \mathfrak{E}' \mathfrak{F}' \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-tensor*:

$\langle \text{kf-filter } (\lambda(x,y). P x \wedge Q y) (\text{kf-tensor } \mathfrak{E} \mathfrak{F}) = \text{kf-tensor } (\text{kf-filter } P \mathfrak{E}) (\text{kf-filter } Q \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-filter-tensor-singleton*:

$\langle \text{kf-filter } ((=) x) (\text{kf-tensor } \mathfrak{E} \mathfrak{F}) = \text{kf-tensor } (\text{kf-filter } ((=) (\text{fst } x)) \mathfrak{E}) (\text{kf-filter } ((=) (\text{snd } x)) \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-cong*:

fixes $\mathfrak{E} \mathfrak{E}' :: \langle ('a \text{ ell2}, 'b \text{ ell2}, 'x) \text{ kraus-family} \rangle$
and $\mathfrak{F} \mathfrak{F}' :: \langle ('c \text{ ell2}, 'd \text{ ell2}, 'y) \text{ kraus-family} \rangle$
assumes $\langle \mathfrak{E} \equiv_{kr} \mathfrak{E}' \rangle$
assumes $\langle \mathfrak{F} \equiv_{kr} \mathfrak{F}' \rangle$
shows $\langle \text{kf-tensor } \mathfrak{E} \mathfrak{F} \equiv_{kr} \text{kf-tensor } \mathfrak{E}' \mathfrak{F}' \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-compose-distrib-weak*:

shows $\langle \text{kf-tensor } (\text{kf-comp } \mathfrak{E} \mathfrak{F}) (\text{kf-comp } \mathfrak{G} \mathfrak{H})$
 $=_{kr} \text{kf-comp } (\text{kf-tensor } \mathfrak{E} \mathfrak{G}) (\text{kf-tensor } \mathfrak{F} \mathfrak{H}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-compose-distrib*:

shows $\langle \text{kf-tensor } (\text{kf-comp } \mathfrak{E} \mathfrak{F}) (\text{kf-comp } \mathfrak{G} \mathfrak{H})$
 $\equiv_{kr} \text{kf-map } (\lambda((e,g),(f,h)). ((e,f),(g,h))) (\text{kf-comp } (\text{kf-tensor } \mathfrak{E} \mathfrak{G}) (\text{kf-tensor } \mathfrak{F} \mathfrak{H})) \rangle$

⟨proof⟩

lemma *kf-tensor-compose-distrib'*:

shows ⟨*kf-comp* (*kf-tensor* \mathfrak{E} \mathfrak{G}) (*kf-tensor* \mathfrak{F} \mathfrak{H})

\equiv_{k_r} *kf-map* ($\lambda((e,f),(g,h)). ((e,g),(f,h))$) (*kf-tensor* (*kf-comp* \mathfrak{E} \mathfrak{F}) (*kf-comp* \mathfrak{G} \mathfrak{H}))⟩

⟨proof⟩

definition *kf-tensor-right* :: ⟨('extra ell2, 'extra ell2) trace-class \Rightarrow ('qu ell2, ('qu \times 'extra) ell2, unit) kraus-family⟩ **where**

— *kf-tensor-right* ϱ maps σ to $\sigma \otimes_o \varrho$

⟨*kf-tensor-right* $\varrho =$ *kf-map-inj* ($\lambda. ()$) (*kf-comp* (*kf-tensor* *kf-id* (*kf-constant-onedim* ϱ)) (*kf-of-op* (*tensor-ell2-right* (*ket* ())))))⟩

definition *kf-tensor-left* :: ⟨('extra ell2, 'extra ell2) trace-class \Rightarrow ('qu ell2, ('extra \times 'qu) ell2, unit) kraus-family⟩ **where**

— *kf-tensor-right* ϱ maps σ to $\varrho \otimes_o \sigma$

⟨*kf-tensor-left* $\varrho =$ *kf-map-inj* ($\lambda. ()$) (*kf-comp* (*kf-tensor* (*kf-constant-onedim* ϱ) *kf-id*) (*kf-of-op* (*tensor-ell2-left* (*ket* ())))))⟩

lemma *kf-apply-tensor-right[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-tensor-right* $\varrho *_{k_r}$ $\sigma =$ *tc-tensor* σ ϱ ⟩

⟨proof⟩

lemma *kf-apply-tensor-left[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-tensor-left* $\varrho *_{k_r}$ $\sigma =$ *tc-tensor* ϱ σ ⟩

⟨proof⟩

lemma *kf-bound-tensor-right[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-bound* (*kf-tensor-right* ϱ) = *norm* $\varrho *_C$ *id-cblinfun*⟩

⟨proof⟩

lemma *kf-bound-tensor-left[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-bound* (*kf-tensor-left* ϱ) = *norm* $\varrho *_C$ *id-cblinfun*⟩

⟨proof⟩

lemma *kf-norm-tensor-right[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-norm* (*kf-tensor-right* ϱ) = *norm* ϱ ⟩

⟨proof⟩

lemma *kf-norm-tensor-left[simp]*:

assumes ⟨ $\varrho \geq 0$ ⟩

shows ⟨*kf-norm* (*kf-tensor-left* ϱ) = *norm* ϱ ⟩

⟨proof⟩

lemma *kf-trace-preserving-tensor*:

assumes $\langle \text{kf-trace-preserving } \mathfrak{E} \rangle$ **and** $\langle \text{kf-trace-preserving } \mathfrak{F} \rangle$
shows $\langle \text{kf-trace-preserving } (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-trace-reducing-tensor*:
assumes $\langle \text{kf-trace-reducing } \mathfrak{E} \rangle$ **and** $\langle \text{kf-trace-reducing } \mathfrak{F} \rangle$
shows $\langle \text{kf-trace-reducing } (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-map-left*:
 $\langle \text{kf-tensor } (\text{kf-map } f \ \mathfrak{E}) \ \mathfrak{F} \equiv_{kr} \text{kf-map } (\text{apfst } f) (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-map-right*:
 $\langle \text{kf-tensor } \mathfrak{E} (\text{kf-map } f \ \mathfrak{F}) \equiv_{kr} \text{kf-map } (\text{apsnd } f) (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-map-both*:
 $\langle \text{kf-tensor } (\text{kf-map } f \ \mathfrak{E}) (\text{kf-map } g \ \mathfrak{F}) \equiv_{kr} \text{kf-map } (\text{map-prod } f \ g) (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-raw-map-inj-both*:
 $\langle \text{kf-tensor-raw } (\text{kf-map-inj } f \ \mathfrak{E}) (\text{kf-map-inj } g \ \mathfrak{F}) = \text{kf-map-inj } (\lambda(E,F,x,y). (E,F,f \ x, g \ y)) (\text{kf-tensor-raw } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-domain-tensor-raw-subset*:
 $\langle \text{kf-domain } (\text{kf-tensor-raw } \mathfrak{E} \ \mathfrak{F}) \subseteq \text{kf-operators } \mathfrak{E} \times \text{kf-operators } \mathfrak{F} \times \text{kf-domain } \mathfrak{E} \times \text{kf-domain } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-tensor-map-inj-both*:
assumes $\langle \text{inj-on } f (\text{kf-domain } \mathfrak{E}) \rangle$
assumes $\langle \text{inj-on } g (\text{kf-domain } \mathfrak{F}) \rangle$
shows $\langle \text{kf-tensor } (\text{kf-map-inj } f \ \mathfrak{E}) (\text{kf-map-inj } g \ \mathfrak{F}) = \text{kf-map-inj } (\text{map-prod } f \ g) (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-tensor-raw*:
shows $\langle \text{kf-operators } (\text{kf-tensor-raw } \mathfrak{E} \ \mathfrak{F}) = \{E \otimes_o F \mid E \ F. E \in \text{kf-operators } \mathfrak{E} \wedge F \in \text{kf-operators } \mathfrak{F}\} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-tensor*:
shows $\langle \text{kf-operators } (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) \subseteq \text{span } \{E \otimes_o F \mid E \ F. E \in \text{kf-operators } \mathfrak{E} \wedge F \in \text{kf-operators } \mathfrak{F}\} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-domain-tensor*: $\langle \text{kf-domain } (\text{kf-tensor } \mathfrak{E} \ \mathfrak{F}) = \text{kf-domain } \mathfrak{E} \times \text{kf-domain } \mathfrak{F} \rangle$

⟨proof⟩

lemma *kf-tensor-raw-0-left[simp]*: ⟨*kf-tensor-raw* 0 $\mathfrak{E} = 0$ ⟩
⟨proof⟩

lemma *kf-tensor-raw-0-right[simp]*: ⟨*kf-tensor-raw* \mathfrak{E} 0 = 0⟩
⟨proof⟩

lemma *kf-tensor-0-left[simp]*: ⟨*kf-tensor* 0 $\mathfrak{E} = 0$ ⟩
⟨proof⟩

lemma *kf-tensor-0-right[simp]*: ⟨*kf-tensor* \mathfrak{E} 0 = 0⟩
⟨proof⟩

lemma *kf-tensor-of-op*:
⟨*kf-tensor* (*kf-of-op* A) (*kf-of-op* B) = *kf-map* ($\lambda(). (((),()))$) (*kf-of-op* (A \otimes_o B))⟩
⟨proof⟩

2.15 Partial trace

definition *kf-partial-trace-right* :: ⟨((*'a* × *'b*) *ell2*, *'a ell2*, *'b*) *kraus-family*⟩ **where**
⟨*kf-partial-trace-right* = *kf-map* ($\lambda(-, b), -$). *inv ket b*)
(*kf-comp* (*kf-of-op* (*tensor-ell2-right* (*ket* ())*))
(*kf-tensor kf-id* (*kf-trace* (*range ket*))))⟩

definition *kf-partial-trace-left* :: ⟨((*'a* × *'b*) *ell2*, *'b ell2*, *'a*) *kraus-family*⟩ **where**
⟨*kf-partial-trace-left* = *kf-map-inj snd* (*kf-comp kf-partial-trace-right* (*kf-of-op swap-ell2*))⟩

lemma *partial-trace-is-kf-partial-trace*:
fixes *t* :: ⟨((*'a* × *'b*) *ell2*, (*'a* × *'b*) *ell2*) *trace-class*⟩
shows ⟨*partial-trace* *t* = *kf-partial-trace-right* $*_{kr}$ *t*⟩
⟨proof⟩

lemma *partial-trace-ignores-kraus-family*:
assumes ⟨*kf-trace-preserving* \mathfrak{E} ⟩
shows ⟨*partial-trace* (*kf-tensor* \mathfrak{F} \mathfrak{E} $*_{kr}$ ρ) = \mathfrak{F} $*_{kr}$ *partial-trace* ρ ⟩
⟨proof⟩

lemma *kf-partial-trace-bound[simp]*:
shows ⟨*kf-bound* *kf-partial-trace-right* = *id-cblinfun*⟩
⟨proof⟩

lemma *kf-partial-trace-norm[simp]*:
shows ⟨*kf-norm* *kf-partial-trace-right* = 1⟩
⟨proof⟩

lemma *kf-partial-trace-right-apply-singleton*:
⟨*kf-partial-trace-right* $*_{kr}$ $@\{x\}$ ρ = *sandwich-tc* (*tensor-ell2-right* (*ket* *x**) ρ)⟩
⟨proof⟩

lemma *kf-partial-trace-left-apply-singleton*:

$\langle \text{kf-partial-trace-left } *_{kr} \text{ @}\{x\} \varrho = \text{sandwich-tc } (\text{tensor-ell2-left } (\text{ket } x) *) \varrho \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-domain-partial-trace-right[simp]*: $\langle \text{kf-domain } \text{kf-partial-trace-right} = \text{UNIV} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-domain-partial-trace-left[simp]*: $\langle \text{kf-domain } \text{kf-partial-trace-left} = \text{UNIV} \rangle$

$\langle \text{proof} \rangle$

2.16 Complete measurement

lemma *complete-measurement-aux*:

fixes B and $F :: \langle 'a::\text{chilbert-space} \Rightarrow_{CL} 'a \times 'a \rangle \text{ set}$

defines $\langle \text{family} \equiv (\lambda x. (\text{selfbutter } (\text{sgn } x), x)) 'B \rangle$

assumes $\langle \text{finite } F \rangle$ and $FB: \langle F \subseteq \text{family} \rangle$ and $\langle \text{is-ortho-set } B \rangle$

shows $\langle (\sum (E, x) \in F. E * o_{CL} E) \leq \text{id-cblinfun} \rangle$

$\langle \text{proof} \rangle$

lemma *complete-measurement-is-kraus-family*:

assumes $\langle \text{is-ortho-set } B \rangle$

shows $\langle \text{kraus-family } ((\lambda x. (\text{selfbutter } (\text{sgn } x), x)) 'B) \rangle$

$\langle \text{proof} \rangle$

lift-definition *kf-complete-measurement* :: $\langle 'a \text{ set} \Rightarrow ('a::\text{chilbert-space}, 'a, 'a) \text{ kraus-family} \rangle$ is

$\langle \lambda B. \text{if } \text{is-ortho-set } B \text{ then } (\lambda x. (\text{selfbutter } (\text{sgn } x), x)) 'B \text{ else } \{\} \rangle$

$\langle \text{proof} \rangle$

definition *kf-complete-measurement-ket* :: $\langle 'a \text{ ell2}, 'a \text{ ell2}, 'a \rangle \text{ kraus-family} \rangle$ where

$\langle \text{kf-complete-measurement-ket} = \text{kf-map-inj } (\text{inv ket}) (\text{kf-complete-measurement } (\text{range ket})) \rangle$

lemma *kf-complete-measurement-domain[simp]*:

assumes $\langle \text{is-ortho-set } B \rangle$

shows $\langle \text{kf-domain } (\text{kf-complete-measurement } B) = B \rangle$

$\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-domain[simp]*:

$\langle \text{kf-domain } \text{kf-complete-measurement-ket} = \text{UNIV} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-kf-map*:

$\langle \text{kf-complete-measurement-ket} \equiv_{kr} \text{kf-map } (\text{inv ket}) (\text{kf-complete-measurement } (\text{range ket})) \rangle$

$\langle \text{proof} \rangle$

lemma *kf-bound-complete-measurement*:

assumes $\langle is-ortho-set\ B \rangle$
shows $\langle kf-bound\ (kf-complete-measurement\ B) \leq id-cblinfun \rangle$
 $\langle proof \rangle$

lemma *kf-norm-complete-measurement*:
assumes $\langle is-ortho-set\ B \rangle$
shows $\langle kf-norm\ (kf-complete-measurement\ B) \leq 1 \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-invalid*:
assumes $\langle \neg is-ortho-set\ B \rangle$
shows $\langle kf-complete-measurement\ B = 0 \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-idem*:
 $\langle kf-comp\ (kf-complete-measurement\ B)\ (kf-complete-measurement\ B)$
 $\equiv_{kr}\ kf-map\ (\lambda b. (b, b))\ (kf-complete-measurement\ B) \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-idem-weak*:
 $\langle kf-comp\ (kf-complete-measurement\ B)\ (kf-complete-measurement\ B)$
 $=_{kr}\ kf-complete-measurement\ B \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-ket-idem*:
 $\langle kf-comp\ kf-complete-measurement-ket\ kf-complete-measurement-ket$
 $\equiv_{kr}\ kf-map\ (\lambda b. (b, b))\ kf-complete-measurement-ket \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-ket-idem-weak*:
 $\langle kf-comp\ kf-complete-measurement-ket\ kf-complete-measurement-ket$
 $=_{kr}\ kf-complete-measurement-ket \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-apply*:
assumes $[simp]: \langle is-ortho-set\ B \rangle$
shows $\langle kf-complete-measurement\ B *_{kr}\ t = (\sum_{\infty x \in B. sandwich-tc\ (selfbutter\ (sgn\ x))\ t) \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-has-sum*:
assumes $\langle is-ortho-set\ B \rangle$
shows $\langle ((\lambda x. sandwich-tc\ (selfbutter\ (sgn\ x))\ \varrho)\ has-sum\ kf-complete-measurement\ B *_{kr}\ \varrho)$
 $B \rangle$
 $\langle proof \rangle$

lemma *kf-complete-measurement-has-sum-onb*:
assumes $\langle is-onb\ B \rangle$

shows $\langle (\lambda x. \text{ sandwich-tc } (\text{selfbutter } x) \varrho) \text{ has-sum kf-complete-measurement } B *_{kr} \varrho \rangle B$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-has-sum*:
 $\langle (\lambda x. \text{ sandwich-tc } (\text{selfbutter } (\text{ket } x)) \varrho) \text{ has-sum kf-complete-measurement-ket } *_{kr} \varrho \rangle UNIV$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-apply-onb*:
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-complete-measurement } B *_{kr} t = (\sum_{\infty x \in B. \text{ sandwich-tc } (\text{selfbutter } x) t) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-apply*: $\langle \text{kf-complete-measurement-ket } *_{kr} t = (\sum_{\infty x. \text{ sandwich-tc } (\text{selfbutter } (\text{ket } x)) t) \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-complete-measurement-onb[simp]*:
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-bound } (\text{kf-complete-measurement } B) = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-bound-complete-measurement-ket[simp]*:
 $\langle \text{kf-bound kf-complete-measurement-ket} = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-complete-measurement-onb[simp]*:
fixes $B :: \langle 'a :: \{\text{not-singleton, chilbert-space}\} \text{ set} \rangle$
assumes $\langle \text{is-onb } B \rangle$
shows $\langle \text{kf-norm } (\text{kf-complete-measurement } B) = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-norm-complete-measurement-ket[simp]*:
 $\langle \text{kf-norm kf-complete-measurement-ket} = 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-diagonal-operator[simp]*:
 $\langle \text{kf-complete-measurement-ket } *_{kr} \text{ diagonal-operator-tc } f = \text{diagonal-operator-tc } f \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-complete-measurement*:
 $\langle \text{kf-operators } (\text{kf-complete-measurement } B) = (\text{selfbutter } o \text{ sgn}) ' B \rangle$ **if** $\langle \text{is-ortho-set } B \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-complete-measurement-invalid*:
 $\langle \text{kf-operators } (\text{kf-complete-measurement } B) = \{\} \rangle$ **if** $\langle \neg \text{is-ortho-set } B \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-operators-complete-measurement-ket*:

⟨*kf-operators kf-complete-measurement-ket* = *range* ($\lambda c.$ *butterfly* (*ket* *c*) (*ket* *c*))⟩
⟨*proof*⟩

lemma *kf-complete-measurement-apply-butterfly*:

assumes ⟨*is-ortho-set* *B*⟩ **and** ⟨*b* ∈ *B*⟩

shows ⟨*kf-complete-measurement* *B* *_{kr} *tc-butterfly* *b* *b* = *tc-butterfly* *b* *b*⟩

⟨*proof*⟩

lemma *kf-complete-measurement-ket-apply-butterfly*:

⟨*kf-complete-measurement-ket* *_{kr} *tc-butterfly* (*ket* *x*) (*ket* *x*) = *tc-butterfly* (*ket* *x*) (*ket* *x*)⟩

⟨*proof*⟩

lemma *kf-map-eq-kf-map-inj-singleton*:

assumes ⟨*card-le-1* (*Rep-kraus-family* \mathfrak{E})⟩

shows ⟨*kf-map* *f* \mathfrak{E} = *kf-map-inj* *f* \mathfrak{E} ⟩

⟨*proof*⟩

lemma *kf-map-eq-kf-map-inj-singleton'*:

assumes ⟨ $\bigwedge y.$ *card-le-1* (*Rep-kraus-family* (*kf-filter* ((=) *y*) \mathfrak{E})⟩

assumes ⟨*inj-on* *f* (*kf-domain* \mathfrak{E})⟩

shows ⟨*kf-map* *f* \mathfrak{E} = *kf-map-inj* *f* \mathfrak{E} ⟩

⟨*proof*⟩

lemma *kf-filter-singleton-kf-complete-measurement*:

assumes ⟨*x* ∈ *B*⟩ **and** ⟨*is-ortho-set* *B*⟩

shows ⟨*kf-filter* ((=) *x*) (*kf-complete-measurement* *B*) = *kf-map-inj* ($\lambda.$ *x*) (*kf-of-op* (*selfbutter* (*sgn* *x*)))⟩

⟨*proof*⟩

lemma *kf-filter-singleton-kf-complete-measurement'*:

assumes ⟨*x* ∈ *B*⟩ **and** ⟨*is-ortho-set* *B*⟩

shows ⟨*kf-filter* ((=) *x*) (*kf-complete-measurement* *B*) = *kf-map* ($\lambda.$ *x*) (*kf-of-op* (*selfbutter* (*sgn* *x*)))⟩

⟨*proof*⟩

lemma *kf-filter-disjoint*:

assumes ⟨ $\bigwedge x.$ *x* ∈ *kf-domain* \mathfrak{E} \implies *P* *x* = *False*⟩

shows ⟨*kf-filter* *P* \mathfrak{E} = 0⟩

⟨*proof*⟩

lemma *kf-complete-measurement-tensor*:

assumes $\langle \text{is-ortho-set } B \rangle$ **and** $\langle \text{is-ortho-set } C \rangle$
shows $\langle \text{kf-map } (\lambda(b,c). b \otimes_s c) \text{ (kf-tensor (kf-complete-measurement } B) \text{ (kf-complete-measurement } C))} \rangle$
 $= \text{kf-complete-measurement } ((\lambda(b,c). b \otimes_s c) \text{ ' } (B \times C)) \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-filter*: $\langle \text{card-le-1 (Rep-kraus-family (kf-filter } P \mathfrak{E}))} \rangle$ **if** $\langle \text{card-le-1 (Rep-kraus-family } \mathfrak{E})} \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-map-inj*[*iff*]: $\langle \text{card-le-1 (Rep-kraus-family (kf-map-inj } f \mathfrak{E}))} \rangle$ **if** $\langle \text{card-le-1 (Rep-kraus-family } \mathfrak{E})} \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-map*[*iff*]: $\langle \text{card-le-1 (Rep-kraus-family (kf-map } f \mathfrak{E}))} \rangle$ **if** $\langle \text{card-le-1 (Rep-kraus-family } \mathfrak{E})} \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-tensor-raw*[*iff*]: $\langle \text{card-le-1 (Rep-kraus-family (kf-tensor-raw } \mathfrak{E} \mathfrak{F}))} \rangle$ **if** $\langle \text{card-le-1 (Rep-kraus-family } \mathfrak{E})} \rangle$ **and** $\langle \text{card-le-1 (Rep-kraus-family } \mathfrak{F})} \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-tensor*[*iff*]: $\langle \text{card-le-1 (Rep-kraus-family (kf-tensor } \mathfrak{E} \mathfrak{F}))} \rangle$ **if** $\langle \text{card-le-1 (Rep-kraus-family } \mathfrak{E})} \rangle$ **and** $\langle \text{card-le-1 (Rep-kraus-family } \mathfrak{F})} \rangle$
 $\langle \text{proof} \rangle$

lemma *card-le-1-kf-filter-complete-measurement*: $\langle \text{card-le-1 (Rep-kraus-family (kf-filter } ((=)x) \text{ (kf-complete-measurement } B))} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-complete-measurement-ket-tensor*:
shows $\langle \text{kf-tensor (kf-complete-measurement-ket :: (-, -, 'a) kraus-family) (kf-complete-measurement-ket :: (-, -, 'b) kraus-family)} \rangle$
 $= \text{kf-complete-measurement-ket} \rangle$
 $\langle \text{proof} \rangle$

2.17 Reconstruction

lemma *kf-reconstruction-is-bounded-clinear*:
assumes $\langle \bigwedge \varrho. ((\lambda a. \text{sandwich-tc } (f \ a) \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ A \rangle$
shows $\langle \text{bounded-clinear } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kf-reconstruction-is-kraus-family*:
assumes *sum*: $\langle \bigwedge \varrho. ((\lambda a. \text{sandwich-tc } (f \ a) \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ A \rangle$
defines $\langle F \equiv \text{Set.filter } (\lambda(E,-). E \neq 0) ((\lambda a. (f \ a, \ a)) \text{ ' } A) \rangle$
shows $\langle \text{kraus-family } F \rangle$

⟨proof⟩

lemma *kf-reconstruction*:

assumes *sum*: ⟨ $\bigwedge \varrho. ((\lambda a. \text{ sandwich-tc } (f \ a) \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ A$ ⟩

defines $\langle F \equiv \text{Abs-kraus-family } (\text{Set.filter } (\lambda(E,-). E \neq 0) ((\lambda a. (f \ a, \ a)) \ 'A)) \rangle$

shows $\langle \text{kf-apply } F = \mathfrak{E} \rangle$

⟨proof⟩

2.18 Cleanup

unbundle *no cblinfun-syntax*

unbundle *no kraus-map-syntax*

end

3 Kraus maps

theory *Kraus-Maps*

imports *Kraus-Families*

begin

3.1 Kraus maps

unbundle *kraus-map-syntax*

unbundle *cblinfun-syntax*

definition *kraus-map* :: $\langle ('a::\text{chilbert-space}, 'a) \text{ trace-class} \Rightarrow ('b::\text{chilbert-space}, 'b) \text{ trace-class} \rangle$
 $\Rightarrow \text{bool}$ **where**

kraus-map-def-raw: $\langle \text{kraus-map } \mathfrak{E} \longleftrightarrow (\exists EE :: ('a, 'b, \text{unit}) \text{ kraus-family}. \mathfrak{E} = \text{kf-apply } EE) \rangle$

lemma *kraus-map-def*: $\langle \text{kraus-map } \mathfrak{E} \longleftrightarrow (\exists EE :: ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x) \text{ kraus-family}. \mathfrak{E} = \text{kf-apply } EE) \rangle$

— Has a more general type than the original definition

⟨proof⟩

lemma *kraus-mapI*:

assumes $\langle \mathfrak{E} = \text{kf-apply } \mathfrak{E}' \rangle$

shows $\langle \text{kraus-map } \mathfrak{E} \rangle$

⟨proof⟩

lemma *kraus-map-bounded-clinear*:

$\langle \text{bounded-clinear } \mathfrak{E} \rangle$ **if** $\langle \text{kraus-map } \mathfrak{E} \rangle$

⟨proof⟩

lemma *kraus-map-pos*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \varrho \geq 0 \rangle$

shows $\langle \mathfrak{E} \ \varrho \geq 0 \rangle$

⟨proof⟩

lemma *kraus-map-mono*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \varrho \geq \tau \rangle$
shows $\langle \mathfrak{E} \varrho \geq \mathfrak{E} \tau \rangle$
\langle proof \rangle

lemma *kraus-map-kf-apply[iff]*: $\langle \text{kraus-map } (\text{kf-apply } \mathfrak{E}) \rangle$
\langle proof \rangle

definition *km-some-kraus-family* :: $\langle ('a::\text{chilbert-space}, 'a) \text{ trace-class} \Rightarrow ('b::\text{chilbert-space}, 'b) \text{ trace-class} \rangle \Rightarrow ('a, 'b, \text{unit}) \text{ kraus-family} \rangle$ **where**
 $\langle \text{km-some-kraus-family } \mathfrak{E} = (\text{if kraus-map } \mathfrak{E} \text{ then SOME } \mathfrak{F}. \mathfrak{E} = \text{kf-apply } \mathfrak{F} \text{ else } 0) \rangle$

lemma *kf-apply-km-some-kraus-family[simp]*:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$
shows $\langle \text{kf-apply } (\text{km-some-kraus-family } \mathfrak{E}) = \mathfrak{E} \rangle$
\langle proof \rangle

lemma *km-some-kraus-family-invalid*:
assumes $\langle \neg \text{kraus-map } \mathfrak{E} \rangle$
shows $\langle \text{km-some-kraus-family } \mathfrak{E} = 0 \rangle$
\langle proof \rangle

definition *km-operators-in* :: $\langle ('a::\text{chilbert-space}, 'a) \text{ trace-class} \Rightarrow ('b::\text{chilbert-space}, 'b) \text{ trace-class} \rangle \Rightarrow ('a \Rightarrow_{CL} 'b) \text{ set} \Rightarrow \text{bool} \rangle$ **where**
 $\langle \text{km-operators-in } \mathfrak{E} S \iff (\exists \mathfrak{F} :: ('a, 'b, \text{unit}) \text{ kraus-family}. \text{kf-apply } \mathfrak{F} = \mathfrak{E} \wedge \text{kf-operators } \mathfrak{F} \subseteq S) \rangle$

lemma *km-operators-in-mono*: $\langle S \subseteq T \implies \text{km-operators-in } \mathfrak{E} S \implies \text{km-operators-in } \mathfrak{E} T \rangle$
\langle proof \rangle

lemma *km-operators-in-kf-apply*:
assumes $\langle \text{span } (\text{kf-operators } \mathfrak{E}) \subseteq S \rangle$
shows $\langle \text{km-operators-in } (\text{kf-apply } \mathfrak{E}) S \rangle$
\langle proof \rangle

lemma *km-operators-in-kf-apply-flattened*:
fixes $\mathfrak{E} :: ('a::\text{chilbert-space}, 'b::\text{chilbert-space}, 'x::\text{CARD-1}) \text{ kraus-family} \rangle$
assumes $\langle \text{kf-operators } \mathfrak{E} \subseteq S \rangle$
shows $\langle \text{km-operators-in } (\text{kf-apply } \mathfrak{E}) S \rangle$
\langle proof \rangle

lemma *km-commute*:
assumes $\langle \text{km-operators-in } \mathfrak{E} S \rangle$
assumes $\langle \text{km-operators-in } \mathfrak{F} T \rangle$
assumes $\langle S \subseteq \text{commutant } T \rangle$
shows $\langle \mathfrak{F} \circ \mathfrak{E} = \mathfrak{E} \circ \mathfrak{F} \rangle$
\langle proof \rangle

lemma *km-operators-in-UNIV*:

assumes $\langle \textit{kraus-map } \mathfrak{E} \rangle$
shows $\langle \textit{km-operators-in } \mathfrak{E} \textit{ UNIV} \rangle$
 $\langle \textit{proof} \rangle$

lemma *separating-kraus-map-bounded-clinear*:
fixes $S :: \langle ('a::\textit{chilbert-space}, 'a) \textit{ trace-class set} \rangle$
assumes $\langle \textit{separating-set (bounded-clinear} :: (- \Rightarrow ('b::\textit{chilbert-space}, 'b) \textit{ trace-class}) \Rightarrow -) S \rangle$
shows $\langle \textit{separating-set (kraus-map} :: (- \Rightarrow ('b::\textit{chilbert-space}, 'b) \textit{ trace-class}) \Rightarrow -) S \rangle$
 $\langle \textit{proof} \rangle$

3.2 Bound and norm

definition *km-bound* :: $\langle (('a::\textit{chilbert-space}, 'a) \textit{ trace-class} \Rightarrow ('b::\textit{chilbert-space}, 'b) \textit{ trace-class}) \Rightarrow ('a, 'a) \textit{ cblinfun} \rangle$ **where**
 $\langle \textit{km-bound } \mathfrak{E} = (\textit{if } \exists \mathfrak{E}' :: (-, -, \textit{unit}) \textit{ kraus-family. } \mathfrak{E} = \textit{kf-apply } \mathfrak{E}' \textit{ then kf-bound (SOME } \mathfrak{E}' :: (-, -, \textit{unit}) \textit{ kraus-family. } \mathfrak{E} = \textit{kf-apply } \mathfrak{E}') \textit{ else } 0) \rangle$

lemma *km-bound-kf-bound*:
assumes $\langle \mathfrak{E} = \textit{kf-apply } \mathfrak{F} \rangle$
shows $\langle \textit{km-bound } \mathfrak{E} = \textit{kf-bound } \mathfrak{F} \rangle$
 $\langle \textit{proof} \rangle$

definition *km-norm* :: $\langle (('a::\textit{chilbert-space}, 'a) \textit{ trace-class} \Rightarrow ('b::\textit{chilbert-space}, 'b) \textit{ trace-class}) \Rightarrow \textit{real} \rangle$ **where**
 $\langle \textit{km-norm } \mathfrak{E} = \textit{norm (km-bound } \mathfrak{E}) \rangle$

lemma *km-norm-kf-norm*:
assumes $\langle \mathfrak{E} = \textit{kf-apply } \mathfrak{F} \rangle$
shows $\langle \textit{km-norm } \mathfrak{E} = \textit{kf-norm } \mathfrak{F} \rangle$
 $\langle \textit{proof} \rangle$

lemma *km-bound-invalid*:
assumes $\langle \neg \textit{kraus-map } \mathfrak{E} \rangle$
shows $\langle \textit{km-bound } \mathfrak{E} = 0 \rangle$
 $\langle \textit{proof} \rangle$

lemma *km-norm-invalid*:
assumes $\langle \neg \textit{kraus-map } \mathfrak{E} \rangle$
shows $\langle \textit{km-norm } \mathfrak{E} = 0 \rangle$
 $\langle \textit{proof} \rangle$

lemma *km-norm-geq0[iff]*: $\langle \textit{km-norm } \mathfrak{E} \geq 0 \rangle$
 $\langle \textit{proof} \rangle$

lemma *kf-bound-pos[iff]*: $\langle \textit{km-bound } \mathfrak{E} \geq 0 \rangle$
 $\langle \textit{proof} \rangle$

lemma *km-bounded-pos*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \varrho \geq 0 \rangle$

shows $\langle \text{norm } (\mathfrak{E} \varrho) \leq \text{km-norm } \mathfrak{E} * \text{norm } \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *km-bounded*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$

shows $\langle \text{norm } (\mathfrak{E} \varrho) \leq 4 * \text{km-norm } \mathfrak{E} * \text{norm } \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-from-map*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$

shows $\langle \psi \cdot_C \text{km-bound } \mathfrak{E} \varphi = \text{trace-tc } (\mathfrak{E} (\text{tc-butterfly } \varphi \psi)) \rangle$

$\langle \text{proof} \rangle$

lemma *trace-from-km-bound*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$

shows $\langle \text{trace-tc } (\mathfrak{E} \varrho) = \text{trace-tc } (\text{compose-tcr } (\text{km-bound } \mathfrak{E}) \varrho) \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-selfadjoint[iff]*: $\langle \text{selfadjoint } (\text{km-bound } \mathfrak{E}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-leq-km-norm-id*: $\langle \text{km-bound } \mathfrak{E} \leq \text{km-norm } \mathfrak{E} *_R \text{id-cblinfun} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-km-some-kraus-family[simp]*: $\langle \text{kf-norm } (\text{km-some-kraus-family } \mathfrak{E}) = \text{km-norm } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

3.3 Basic Kraus maps

Zero map and constant maps. Addition and rescaling and composition of maps.

lemma *kraus-map-0[iff]*: $\langle \text{kraus-map } 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-0'[iff]*: $\langle \text{kraus-map } (\lambda-. 0) \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-0[simp]*: $\langle \text{km-bound } 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *km-norm-0[simp]*: $\langle \text{km-norm } 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *km-some-kraus-family-0[simp]*: $\langle \text{km-some-kraus-family } 0 = 0 \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-id*[*iff*]: $\langle \text{kraus-map id} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-id*[*simp*]: $\langle \text{km-bound id} = \text{id-cblinfun} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-norm-id-leq1*[*iff*]: $\langle \text{km-norm id} \leq 1 \rangle$
 $\langle \text{proof} \rangle$

lemma *km-norm-id-eq1*[*simp*]: $\langle \text{km-norm (id :: ('a :: \{chilbert-space, not-singleton\}, 'a) \text{ trace-class} \Rightarrow \cdot) = 1} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-operators-in-id*[*iff*]: $\langle \text{km-operators-in id} \{ \text{id-cblinfun} \} \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-add*[*iff*]:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{kraus-map } (\lambda \varrho. \mathfrak{E} \varrho + \mathfrak{F} \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-plus'*[*iff*]:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{kraus-map } (\mathfrak{E} + \mathfrak{F}) \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-plus*:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{km-bound } (\mathfrak{E} + \mathfrak{F}) = \text{km-bound } \mathfrak{E} + \text{km-bound } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-norm-triangle*:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{km-norm } (\mathfrak{E} + \mathfrak{F}) \leq \text{km-norm } \mathfrak{E} + \text{km-norm } \mathfrak{F} \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-constant*[*iff*]: $\langle \text{kraus-map } (\lambda \sigma. \text{trace-tc } \sigma *_C \varrho) \rangle$ **if** $\langle \varrho \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-constant-invalid*:
 $\langle \neg \text{kraus-map } (\lambda \sigma :: ('a :: \{chilbert-space, not-singleton\}, 'a) \text{ trace-class. trace-tc } \sigma *_C \varrho) \rangle$ **if** $\langle \sim \varrho \geq 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-scale*:
assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle c \geq 0 \rangle$
shows $\langle \text{kraus-map } (\lambda \varrho. c *_R \mathfrak{E} \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-scale[simp]*: $\langle km\text{-bound } (\lambda \varrho. c *_{\mathbb{R}} \mathfrak{E} \varrho) = c *_{\mathbb{R}} km\text{-bound } \mathfrak{E} \rangle$ **if** $\langle c \geq 0 \rangle$
 $\langle proof \rangle$

lemma *km-norm-scale[simp]*: $\langle km\text{-norm } (\lambda \varrho. c *_{\mathbb{R}} \mathfrak{E} \varrho) = c * km\text{-norm } \mathfrak{E} \rangle$ **if** $\langle c \geq 0 \rangle$
 $\langle proof \rangle$

lemma *kraus-map-sandwich[iff]*: $\langle kraus\text{-map } (sandwich\text{-tc } A) \rangle$
 $\langle proof \rangle$

lemma *km-bound-sandwich[simp]*: $\langle km\text{-bound } (sandwich\text{-tc } A) = A * o_{CL} A \rangle$
 $\langle proof \rangle$

lemma *km-norm-sandwich[simp]*: $\langle km\text{-norm } (sandwich\text{-tc } A) = (norm A)^2 \rangle$
 $\langle proof \rangle$

lemma *km-operators-in-sandwich*: $\langle km\text{-operators-in } (sandwich\text{-tc } U) \{U\} \rangle$
 $\langle proof \rangle$

lemma *km-constant-bound[simp]*: $\langle km\text{-bound } (\lambda \sigma. trace\text{-tc } \sigma *_{\mathbb{C}} \varrho) = norm \varrho *_{\mathbb{R}} id\text{-cblinfun} \rangle$ **if**
 $\langle \varrho \geq 0 \rangle$
 $\langle proof \rangle$

lemma *km-constant-norm[simp]*: $\langle km\text{-norm } (\lambda \sigma :: ('a :: \{chilbert\text{-space}, not\text{-singleton}\}, 'a) trace\text{-class. } trace\text{-tc } \sigma *_{\mathbb{C}} \varrho) = norm \varrho \rangle$ **if** $\langle \varrho \geq 0 \rangle$
 $\langle proof \rangle$

lemma *km-constant-norm-leq[simp]*: $\langle km\text{-norm } (\lambda \sigma :: ('a :: chilbert\text{-space}, 'a) trace\text{-class. } trace\text{-tc } \sigma *_{\mathbb{C}} \varrho) \leq norm \varrho \rangle$
 $\langle proof \rangle$

lemma *kraus-map-comp*:
assumes $\langle kraus\text{-map } \mathfrak{E} \rangle$ **and** $\langle kraus\text{-map } \mathfrak{F} \rangle$
shows $\langle kraus\text{-map } (\mathfrak{E} \circ \mathfrak{F}) \rangle$
 $\langle proof \rangle$

lemma *km-comp-norm-leq*:
assumes $\langle kraus\text{-map } \mathfrak{E} \rangle$ **and** $\langle kraus\text{-map } \mathfrak{F} \rangle$
shows $\langle km\text{-norm } (\mathfrak{E} \circ \mathfrak{F}) \leq km\text{-norm } \mathfrak{E} * km\text{-norm } \mathfrak{F} \rangle$
 $\langle proof \rangle$

lemma *km-bound-comp-sandwich*:
assumes $\langle kraus\text{-map } \mathfrak{E} \rangle$
shows $\langle km\text{-bound } (\lambda \varrho. \mathfrak{E} (sandwich\text{-tc } U \varrho)) = sandwich (U*) (km\text{-bound } \mathfrak{E}) \rangle$
 $\langle proof \rangle$

lemma *km-norm-comp-sandwich-coiso*:
assumes $\langle \text{isometry } (U^*) \rangle$
shows $\langle \text{km-norm } (\lambda \varrho. \mathfrak{E} (\text{sandwich-tc } U \ \varrho)) = \text{km-norm } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-comp-sandwich-iso*:
assumes $\langle \text{isometry } U \rangle$
shows $\langle \text{km-bound } (\lambda \varrho. \text{sandwich-tc } U \ (\mathfrak{E} \ \varrho)) = \text{km-bound } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *km-norm-comp-sandwich-iso*:
assumes $\langle \text{isometry } U \rangle$
shows $\langle \text{km-norm } (\lambda \varrho. \text{sandwich-tc } U \ (\mathfrak{E} \ \varrho)) = \text{km-norm } \mathfrak{E} \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-sum*:
assumes $\langle \bigwedge x. x \in A \implies \text{kraus-map } (\mathfrak{E} \ x) \rangle$
shows $\langle \text{kraus-map } (\sum x \in A. \mathfrak{E} \ x) \rangle$
 $\langle \text{proof} \rangle$

lemma *km-bound-sum*:
assumes $\langle \bigwedge x. x \in A \implies \text{kraus-map } (\mathfrak{E} \ x) \rangle$
shows $\langle \text{km-bound } (\sum x \in A. \mathfrak{E} \ x) = (\sum x \in A. \text{km-bound } (\mathfrak{E} \ x)) \rangle$
 $\langle \text{proof} \rangle$

3.4 Infinite sums

lemma
assumes $\langle \bigwedge \varrho. ((\lambda a. \text{sandwich-tc } (f \ a) \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ A \rangle$
defines $\langle EE \equiv \text{Set.filter } (\lambda (E, -). E \neq 0) ((\lambda a. (f \ a, \ a)) \text{ ' } A) \rangle$
shows *kraus-mapI-sum*: $\langle \text{kraus-map } \mathfrak{E} \rangle$
and *kraus-map-sum-kraus-family*: $\langle \text{kraus-family } EE \rangle$
and *kraus-map-sum-kf-apply*: $\langle \mathfrak{E} = \text{kf-apply } (\text{Abs-kraus-family } EE) \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-infsum-sandwich*:
assumes $\langle \bigwedge \varrho. (\lambda a. \text{sandwich-tc } (f \ a) \ \varrho) \text{ summable-on } A \rangle$
shows $\langle \text{kraus-map } (\lambda \varrho. \sum_{\infty} a \in A. \text{sandwich-tc } (f \ a) \ \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-sum-sandwich*: $\langle \text{kraus-map } (\lambda \varrho. \sum a \in A. \text{sandwich-tc } (f \ a) \ \varrho) \rangle$
 $\langle \text{proof} \rangle$

lemma *kraus-map-as-infsum*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$

shows $\langle \exists M. \forall \varrho. ((\lambda E. \text{sandwich-tc } E \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ M \rangle$

$\langle \text{proof} \rangle$

definition *km-summable* :: $\langle ('a \Rightarrow ('b::\text{chilbert-space}, 'b) \text{ trace-class} \Rightarrow ('c::\text{chilbert-space}, 'c) \text{ trace-class}) \Rightarrow 'a \text{ set} \Rightarrow \text{bool} \rangle$ **where**

$\langle \text{km-summable } \mathfrak{E} \ A \iff \text{summable-on-in cweak-operator-topology } (\lambda a. \text{km-bound } (\mathfrak{E} \ a)) \ A \rangle$

lemma *km-summable-kf-summable*:

assumes $\langle \bigwedge a \ \varrho. a \in A \implies \mathfrak{E} \ a \ \varrho = \mathfrak{F} \ a \ *_{kr} \ \varrho \rangle$

shows $\langle \text{km-summable } \mathfrak{E} \ A \iff \text{kf-summable } \mathfrak{F} \ A \rangle$

$\langle \text{proof} \rangle$

lemma *km-summable-summable*:

assumes *km*: $\langle \bigwedge a. a \in A \implies \text{kraus-map } (\mathfrak{E} \ a) \rangle$

assumes *sum*: $\langle \text{km-summable } \mathfrak{E} \ A \rangle$

shows $\langle (\lambda a. \mathfrak{E} \ a \ \varrho) \text{ summable-on } A \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-infsum*:

assumes *km*: $\langle \bigwedge a. a \in A \implies \text{kraus-map } (\mathfrak{E} \ a) \rangle$

assumes *sum*: $\langle \text{km-summable } \mathfrak{E} \ A \rangle$

shows $\langle \text{kraus-map } (\lambda \varrho. \sum_{\infty a \in A. \mathfrak{E} \ a \ \varrho}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-infsum*:

assumes *km*: $\langle \bigwedge a. a \in A \implies \text{kraus-map } (\mathfrak{E} \ a) \rangle$

assumes *sum*: $\langle \text{km-summable } \mathfrak{E} \ A \rangle$

shows $\langle \text{km-bound } (\lambda \varrho. \sum_{\infty a \in A. \mathfrak{E} \ a \ \varrho}) = \text{infsum-in cweak-operator-topology } (\lambda a. \text{km-bound } (\mathfrak{E} \ a)) \ A \rangle$

$\langle \text{proof} \rangle$

lemma *km-norm-infsum*:

assumes *km*: $\langle \bigwedge a. a \in A \implies \text{kraus-map } (\mathfrak{E} \ a) \rangle$

assumes *sum*: $\langle (\lambda a. \text{km-norm } (\mathfrak{E} \ a)) \text{ summable-on } A \rangle$

shows $\langle \text{km-norm } (\lambda \varrho. \sum_{\infty a \in A. \mathfrak{E} \ a \ \varrho}) \leq (\sum_{\infty a \in A. \text{km-norm } (\mathfrak{E} \ a)) \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-has-sum*:

assumes $\langle \bigwedge x. x \in A \implies \text{kraus-map } (\mathfrak{E} \ x) \rangle$

assumes $\langle \text{km-summable } \mathfrak{E} \ A \rangle$

assumes $\langle (\mathfrak{E} \ \text{has-sum } \mathfrak{F}) \ A \rangle$

shows $\langle \text{kraus-map } \mathfrak{F} \rangle$

<proof>

lemma *km-summable-iff-sums-to-kraus-map*:

assumes $\langle \bigwedge a. a \in A \implies \text{kraus-map } (\mathfrak{E} a) \rangle$

shows $\langle \text{km-summable } \mathfrak{E} A \longleftrightarrow (\exists \mathfrak{F}. (\forall t. ((\lambda x. \mathfrak{E} x t) \text{ has-sum } \mathfrak{F} t) A) \wedge \text{kraus-map } \mathfrak{F}) \rangle$

<proof>

3.5 Tensor products

definition *km-tensor-exists* :: $\langle ((\text{'a ell2}, \text{'b ell2}) \text{ trace-class} \implies (\text{'c ell2}, \text{'d ell2}) \text{ trace-class})$
 $\implies ((\text{'e ell2}, \text{'f ell2}) \text{ trace-class} \implies (\text{'g ell2}, \text{'h ell2}) \text{ trace-class}) \implies \text{bool} \rangle$

where

$\langle \text{km-tensor-exists } \mathfrak{E} \mathfrak{F} \longleftrightarrow (\exists \mathfrak{E}\mathfrak{F}. \text{bounded-clinear } \mathfrak{E}\mathfrak{F} \wedge (\forall \varrho \sigma. \mathfrak{E}\mathfrak{F} (\text{tc-tensor } \varrho \sigma) = \text{tc-tensor } (\mathfrak{E} \varrho) (\mathfrak{F} \sigma))) \rangle$

definition *km-tensor* :: $\langle ((\text{'a ell2}, \text{'c ell2}) \text{ trace-class} \implies (\text{'e ell2}, \text{'g ell2}) \text{ trace-class})$
 $\implies ((\text{'b ell2}, \text{'d ell2}) \text{ trace-class} \implies (\text{'f ell2}, \text{'h ell2}) \text{ trace-class})$
 $\implies ((\text{'a} \times \text{'b}) \text{ ell2}, (\text{'c} \times \text{'d}) \text{ ell2}) \text{ trace-class} \implies ((\text{'e} \times \text{'f}) \text{ ell2}, (\text{'g} \times \text{'h}) \text{ ell2})$

trace-class **where**

$\langle \text{km-tensor } \mathfrak{E} \mathfrak{F} = (\text{if } \text{km-tensor-exists } \mathfrak{E} \mathfrak{F}$
 $\text{ then SOME } \mathfrak{E}\mathfrak{F}. \text{bounded-clinear } \mathfrak{E}\mathfrak{F} \wedge (\forall \varrho \sigma. \mathfrak{E}\mathfrak{F} (\text{tc-tensor } \varrho \sigma) = \text{tc-tensor } (\mathfrak{E} \varrho) (\mathfrak{F} \sigma))$
 $\text{ else } 0) \rangle$

lemma *km-tensor-invalid*:

assumes $\langle \neg \text{km-tensor-exists } \mathfrak{E} \mathfrak{F} \rangle$

shows $\langle \text{km-tensor } \mathfrak{E} \mathfrak{F} = 0 \rangle$

<proof>

lemma *km-tensor-exists-bounded-clinear*[*iff*]:

assumes $\langle \text{km-tensor-exists } \mathfrak{E} \mathfrak{F} \rangle$

shows $\langle \text{bounded-clinear } (\text{km-tensor } \mathfrak{E} \mathfrak{F}) \rangle$

<proof>

lemma *km-tensor-apply*[*simp*]:

assumes $\langle \text{km-tensor-exists } \mathfrak{E} \mathfrak{F} \rangle$

shows $\langle \text{km-tensor } \mathfrak{E} \mathfrak{F} (\text{tc-tensor } \varrho \sigma) = \text{tc-tensor } (\mathfrak{E} \varrho) (\mathfrak{F} \sigma) \rangle$

<proof>

lemma *km-tensor-unique*:

assumes $\langle \text{bounded-clinear } \mathfrak{E}\mathfrak{F} \rangle$

assumes $\langle \bigwedge \varrho \sigma. \mathfrak{E}\mathfrak{F} (\text{tc-tensor } \varrho \sigma) = \text{tc-tensor } (\mathfrak{E} \varrho) (\mathfrak{F} \sigma) \rangle$

shows $\langle \mathfrak{E}\mathfrak{F} = \text{km-tensor } \mathfrak{E} \mathfrak{F} \rangle$

<proof>

lemma *km-tensor-kf-tensor*: $\langle \text{km-tensor } (\text{kf-apply } \mathfrak{E}) (\text{kf-apply } \mathfrak{F}) = \text{kf-apply } (\text{km-tensor } \mathfrak{E} \mathfrak{F}) \rangle$

<proof>

lemma *km-tensor-kraus-map*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{kraus-map } (\text{km-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-kraus-map-exists*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{km-tensor-exists } \mathfrak{E} \ \mathfrak{F} \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-as-infsum*:

assumes $\langle \bigwedge \varrho. ((\lambda i. \text{sandwich-tc } (E \ i) \ \varrho) \text{ has-sum } \mathfrak{E} \ \varrho) \ I \rangle$
assumes $\langle \bigwedge \varrho. ((\lambda j. \text{sandwich-tc } (F \ j) \ \varrho) \text{ has-sum } \mathfrak{F} \ \varrho) \ J \rangle$
shows $\langle \text{km-tensor } \mathfrak{E} \ \mathfrak{F} \ \varrho = (\sum_{\infty (i,j) \in I \times J. \text{sandwich-tc } (E \ i \otimes_o F \ j) \ \varrho} \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-tensor*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{km-bound } (\text{km-tensor } \mathfrak{E} \ \mathfrak{F}) = \text{km-bound } \mathfrak{E} \otimes_o \text{km-bound } \mathfrak{F} \rangle$

$\langle \text{proof} \rangle$

lemma *km-norm-tensor*:

assumes $\langle \text{kraus-map } \mathfrak{E} \rangle$ **and** $\langle \text{kraus-map } \mathfrak{F} \rangle$
shows $\langle \text{km-norm } (\text{km-tensor } \mathfrak{E} \ \mathfrak{F}) = \text{km-norm } \mathfrak{E} * \text{km-norm } \mathfrak{F} \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-compose-distrib*:

assumes $\langle \text{km-tensor-exists } \mathfrak{E} \ \mathfrak{G} \rangle$ **and** $\langle \text{km-tensor-exists } \mathfrak{F} \ \mathfrak{H} \rangle$
shows $\langle \text{km-tensor } (\mathfrak{E} \ o \ \mathfrak{F}) \ (\mathfrak{G} \ o \ \mathfrak{H}) = \text{km-tensor } \mathfrak{E} \ \mathfrak{G} \ o \ \text{km-tensor } \mathfrak{F} \ \mathfrak{H} \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-tensor-right[simp]*:

assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kraus-map } (\lambda \sigma. \text{tc-tensor } \sigma \ \varrho) \rangle$

$\langle \text{proof} \rangle$

lemma *kraus-map-tensor-left[simp]*:

assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{kraus-map } (\lambda \sigma. \text{tc-tensor } \varrho \ \sigma) \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-tensor-right[simp]*:

assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{km-bound } (\lambda \sigma. \text{tc-tensor } \sigma \ \varrho) = \text{norm } \varrho *_{\mathcal{C}} \text{id-cblinfun} \rangle$

$\langle \text{proof} \rangle$

lemma *km-bound-tensor-left[simp]*:

assumes $\langle \varrho \geq 0 \rangle$
shows $\langle \text{km-bound } (\lambda \sigma. \text{tc-tensor } \varrho \ \sigma) = \text{norm } \varrho *_{\mathcal{C}} \text{id-cblinfun} \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-tensor-right[simp]*:

assumes $\langle \varrho \geq 0 \rangle$

shows $\langle \text{km-norm } (\lambda\sigma. \text{tc-tensor } \sigma \ \varrho) = \text{norm } \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *kf-norm-tensor-left[simp]*:

assumes $\langle \varrho \geq 0 \rangle$

shows $\langle \text{km-norm } (\lambda\sigma. \text{tc-tensor } \varrho \ \sigma) = \text{norm } \varrho \rangle$

$\langle \text{proof} \rangle$

lemma *km-operators-in-tensor*:

assumes $\langle \text{km-operators-in } \mathfrak{E} \ S \rangle$

assumes $\langle \text{km-operators-in } \mathfrak{F} \ T \rangle$

shows $\langle \text{km-operators-in } (\text{km-tensor } \mathfrak{E} \ \mathfrak{F}) \ (\text{span } \{s \otimes_o t \mid s \in S \wedge t \in T\}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-tensor-sandwich-tc*:

$\langle \text{km-tensor } (\text{sandwich-tc } A) \ (\text{sandwich-tc } B) = \text{sandwich-tc } (A \otimes_o B) \rangle$

$\langle \text{proof} \rangle$

3.6 Trace and partial trace

definition $\langle \text{km-trace-preserving } \mathfrak{E} \longleftrightarrow (\exists \mathfrak{F}::(-, -, \text{unit}) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } \mathfrak{F} \wedge \text{kf-trace-preserving } \mathfrak{F}) \rangle$

lemma *km-trace-preserving-def'*: $\langle \text{km-trace-preserving } \mathfrak{E} \longleftrightarrow (\exists \mathfrak{F}::(-, -, 'c) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } \mathfrak{F} \wedge \text{kf-trace-preserving } \mathfrak{F}) \rangle$

— Has a more general type than *km-trace-preserving-def*

$\langle \text{proof} \rangle$

definition *km-trace-reducing-def*: $\langle \text{km-trace-reducing } \mathfrak{E} \longleftrightarrow (\exists \mathfrak{F}::(-, -, \text{unit}) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } \mathfrak{F} \wedge \text{kf-trace-reducing } \mathfrak{F}) \rangle$

lemma *km-trace-reducing-def'*: $\langle \text{km-trace-reducing } \mathfrak{E} \longleftrightarrow (\exists \mathfrak{F}::(-, -, 'c) \text{ kraus-family. } \mathfrak{E} = \text{kf-apply } \mathfrak{F} \wedge \text{kf-trace-reducing } \mathfrak{F}) \rangle$

$\langle \text{proof} \rangle$

lemma *km-trace-preserving-apply[simp]*: $\langle \text{km-trace-preserving } (\text{kf-apply } \mathfrak{E}) = \text{kf-trace-preserving } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *km-trace-reducing-apply[simp]*: $\langle \text{km-trace-reducing } (\text{kf-apply } \mathfrak{E}) = \text{kf-trace-reducing } \mathfrak{E} \rangle$

$\langle \text{proof} \rangle$

lemma *km-trace-preserving-iff*: $\langle \text{km-trace-preserving } \mathfrak{E} \longleftrightarrow \text{kraus-map } \mathfrak{E} \wedge (\forall \varrho. \text{trace-tc } (\mathfrak{E} \ \varrho) = \text{trace-tc } \varrho) \rangle$

$\langle \text{proof} \rangle$

lemma *km-trace-reducing-iff*: $\langle km\text{-trace-reducing } \mathfrak{E} \longleftrightarrow kraus\text{-map } \mathfrak{E} \wedge (\forall \varrho \geq 0. trace\text{-tc } (\mathfrak{E} \ \varrho) \leq trace\text{-tc } \varrho) \rangle$
 $\langle proof \rangle$

lemma *km-trace-preserving-imp-reducing*:
assumes $\langle km\text{-trace-preserving } \mathfrak{E} \rangle$
shows $\langle km\text{-trace-reducing } \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *km-trace-preserving-id[iff]*: $\langle km\text{-trace-preserving } id \rangle$
 $\langle proof \rangle$

lemma *km-trace-reducing-iff-norm-leq1*: $\langle km\text{-trace-reducing } \mathfrak{E} \longleftrightarrow kraus\text{-map } \mathfrak{E} \wedge km\text{-norm } \mathfrak{E} \leq 1 \rangle$
 $\langle proof \rangle$

lemma *km-trace-preserving-iff-bound-id*: $\langle km\text{-trace-preserving } \mathfrak{E} \longleftrightarrow kraus\text{-map } \mathfrak{E} \wedge km\text{-bound } \mathfrak{E} = id\text{-cblinfun} \rangle$
 $\langle proof \rangle$

lemma *km-trace-preserving-iff-bound-id'*:
fixes $\mathfrak{E} :: \langle 'a :: \{chilbert\text{-space}, not\text{-singleton}\}, 'a \rangle trace\text{-class} \Rightarrow - \rangle$
shows $\langle km\text{-trace-preserving } \mathfrak{E} \longleftrightarrow km\text{-bound } \mathfrak{E} = id\text{-cblinfun} \rangle$
 $\langle proof \rangle$

lemma *km-trace-norm-preserving*: $\langle km\text{-norm } \mathfrak{E} \leq 1 \rangle$ **if** $\langle km\text{-trace-preserving } \mathfrak{E} \rangle$
 $\langle proof \rangle$

lemma *km-trace-norm-preserving-eq*:
fixes $\mathfrak{E} :: \langle 'a :: \{chilbert\text{-space}, not\text{-singleton}\}, 'a \rangle trace\text{-class} \Rightarrow ('b :: chilbert\text{-space}, 'b) trace\text{-class} \rangle$
assumes $\langle km\text{-trace-preserving } \mathfrak{E} \rangle$
shows $\langle km\text{-norm } \mathfrak{E} = 1 \rangle$
 $\langle proof \rangle$

lemma *kraus-map-trace*: $\langle kraus\text{-map } (one\text{-dim-iso } o \ trace\text{-tc}) \rangle$
 $\langle proof \rangle$

lemma *trace-preserving-trace-kraus-map[iff]*: $\langle km\text{-trace-preserving } (one\text{-dim-iso } o \ trace\text{-tc}) \rangle$
 $\langle proof \rangle$

lemma *km-trace-bound[simp]*: $\langle km\text{-bound } (one\text{-dim-iso } o \ trace\text{-tc}) = id\text{-cblinfun} \rangle$
 $\langle proof \rangle$

lemma *km-trace-norm-eq1[simp]*: $\langle km\text{-norm } (one\text{-dim-iso } o \ trace\text{-tc} :: ('a :: \{chilbert\text{-space}, not\text{-singleton}\}, 'a) trace\text{-class} \Rightarrow -) = 1 \rangle$
 $\langle proof \rangle$

lemma *km-trace-norm-leq1*[simp]: $\langle km\text{-norm } (one\text{-dim-iso } o \text{ trace-}tc) \leq 1 \rangle$
 $\langle proof \rangle$

lemma *kraus-map-partial-trace*[iff]: $\langle kraus\text{-map } partial\text{-trace} \rangle$
 $\langle proof \rangle$

lemma *partial-trace-ignores-kraus-map*:
assumes $\langle km\text{-trace-preserving } \mathfrak{E} \rangle$
assumes $\langle kraus\text{-map } \mathfrak{F} \rangle$
shows $\langle partial\text{-trace } (km\text{-tensor } \mathfrak{F} \ \mathfrak{E} \ \varrho) = \mathfrak{F} \ (partial\text{-trace } \varrho) \rangle$
 $\langle proof \rangle$

lemma *km-partial-trace-bound*[simp]: $\langle km\text{-bound } partial\text{-trace} = id\text{-cblinfun} \rangle$
 $\langle proof \rangle$

lemma *km-partial-trace-norm*[simp]:
shows $\langle km\text{-norm } partial\text{-trace} = 1 \rangle$
 $\langle proof \rangle$

lemma *km-trace-preserving-tensor*:
assumes $\langle km\text{-trace-preserving } \mathfrak{E} \rangle$ **and** $\langle km\text{-trace-preserving } \mathfrak{F} \rangle$
shows $\langle km\text{-trace-preserving } (km\text{-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle proof \rangle$

lemma *km-trace-reducing-tensor*:
assumes $\langle km\text{-trace-reducing } \mathfrak{E} \rangle$ **and** $\langle km\text{-trace-reducing } \mathfrak{F} \rangle$
shows $\langle km\text{-trace-reducing } (km\text{-tensor } \mathfrak{E} \ \mathfrak{F}) \rangle$
 $\langle proof \rangle$

3.7 Complete measurements

definition $\langle km\text{-complete-measurement } B \ \varrho = (\sum_{x \in B} sandwich\text{-}tc \ (selfbutter \ (sgn \ x)) \ \varrho) \rangle$
abbreviation $\langle km\text{-complete-measurement-ket} \equiv km\text{-complete-measurement } (range \ ket) \rangle$

lemma *km-complete-measurement-kf-complete-measurement*: $\langle km\text{-complete-measurement } B = kf\text{-apply } (kf\text{-complete-measurement } B) \rangle$ **if** $\langle is\text{-ortho-set } B \rangle$
 $\langle proof \rangle$

lemma *km-complete-measurement-ket-kf-complete-measurement-ket*: $\langle km\text{-complete-measurement-ket} = kf\text{-apply } kf\text{-complete-measurement-ket} \rangle$
 $\langle proof \rangle$

lemma *km-complete-measurement-has-sum*:
assumes $\langle is\text{-ortho-set } B \rangle$
shows $\langle ((\lambda x. sandwich\text{-}tc \ (selfbutter \ (sgn \ x)) \ \varrho) \text{ has-sum } km\text{-complete-measurement } B \ \varrho) \ B \rangle$

⟨proof⟩

lemma *km-complete-measurement-ket-has-sum*:

⟨((λx. sandwich-tc (selfbutter (ket x)) ρ) has-sum km-complete-measurement-ket ρ) UNIV⟩

⟨proof⟩

lemma *km-bound-complete-measurement*:

assumes ⟨is-ortho-set B⟩

shows ⟨km-bound (km-complete-measurement B) ≤ id-cblinfun⟩

⟨proof⟩

lemma *km-norm-complete-measurement*:

assumes ⟨is-ortho-set B⟩

shows ⟨km-norm (km-complete-measurement B) ≤ 1⟩

⟨proof⟩

lemma *km-bound-complete-measurement-onb[simp]*:

assumes ⟨is-onb B⟩

shows ⟨km-bound (km-complete-measurement B) = id-cblinfun⟩

⟨proof⟩

lemma *km-bound-complete-measurement-ket[simp]*: ⟨km-bound km-complete-measurement-ket = id-cblinfun⟩

⟨proof⟩

lemma *km-norm-complete-measurement-onb[simp]*:

fixes B :: ⟨'a::{not-singleton, chilbert-space} set⟩

assumes ⟨is-onb B⟩

shows ⟨km-norm (km-complete-measurement B) = 1⟩

⟨proof⟩

lemma *km-norm-complete-measurement-ket[simp]*:

shows ⟨km-norm km-complete-measurement-ket = 1⟩

⟨proof⟩

lemma *kraus-map-complete-measurement*:

assumes ⟨is-ortho-set B⟩

shows ⟨kraus-map (km-complete-measurement B)⟩

⟨proof⟩

lemma *kraus-map-complete-measurement-ket[iff]*:

shows ⟨kraus-map km-complete-measurement-ket⟩

⟨proof⟩

lemma *km-complete-measurement-idem[simp]*:

assumes ⟨is-ortho-set B⟩

shows ⟨km-complete-measurement B (km-complete-measurement B ρ) = km-complete-measurement B ρ⟩

⟨proof⟩

lemma *km-complete-measurement-ket-idem*[simp]:
 $\langle km\text{-complete-measurement-ket } (km\text{-complete-measurement-ket } \varrho) = km\text{-complete-measurement-ket } \varrho \rangle$
 ⟨proof⟩

lemma *km-complete-measurement-has-sum-onb*:
assumes ⟨is-onb B⟩
shows ⟨((λx. sandwich-tc (selfbutter x) ρ) has-sum km-complete-measurement B ρ) B⟩
 ⟨proof⟩

lemma *km-complete-measurement-ket-diagonal-operator*[simp]:
 $\langle km\text{-complete-measurement-ket } (diagonal\text{-operator-tc } f) = diagonal\text{-operator-tc } f \rangle$
 ⟨proof⟩

lemma *km-operators-complete-measurement*:
assumes ⟨is-ortho-set B⟩
shows ⟨km-operators-in (km-complete-measurement B) (span (selfbutter ‘ B))⟩
 ⟨proof⟩

lemma *km-operators-complete-measurement-ket*:
shows ⟨km-operators-in km-complete-measurement-ket (span (range (λc. (selfbutter (ket c))))))⟩
 ⟨proof⟩

lemma *km-complete-measurement-ket-butterket*[simp]:
 $\langle km\text{-complete-measurement-ket } (tc\text{-butterfly } (ket\ c) (ket\ c)) = tc\text{-butterfly } (ket\ c) (ket\ c) \rangle$
 ⟨proof⟩

lemma *km-complete-measurement-tensor*:
assumes ⟨is-ortho-set B⟩ **and** ⟨is-ortho-set C⟩
shows ⟨km-tensor (km-complete-measurement B) (km-complete-measurement C)
 = km-complete-measurement ((λ(b,c). b ⊗_s c) ‘ (B × C))⟩
 ⟨proof⟩

lemma *km-complete-measurement-ket-tensor*:
shows ⟨km-tensor (km-complete-measurement-ket :: ('a ell2, -) trace-class ⇒ -) (km-complete-measurement-ket
 :: ('b ell2, -) trace-class ⇒ -)
 = km-complete-measurement-ket⟩
 ⟨proof⟩

lemma *km-tensor-0-left*[simp]: ⟨km-tensor (0 :: ('a ell2, 'b ell2) trace-class ⇒ ('c ell2, 'd ell2)
 trace-class) ℰ = 0⟩
 ⟨proof⟩

lemma *km-tensor-0-right*[simp]: ⟨km-tensor ℰ (0 :: ('a ell2, 'b ell2) trace-class ⇒ ('c ell2, 'd
 ell2) trace-class) = 0⟩
 ⟨proof⟩

unbundle *no kraus-map-syntax*
unbundle *no cblinfun-syntax*

end

References

- [1] K. Kraus. *States, effects, and operations: fundamental notions of quantum theory*, volume 190 of *LNP*. Springer, 1983. Defines Kraus maps in Section 3, Theorem 1. Only considers separable Hilbert spaces.