

The Kolmogorov-Chentsov Theorem

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Abstract

Continuous-time stochastic processes often carry the condition of having almost-surely continuous paths. If some process X satisfies certain bounds on its expectation, then the Kolmogorov-Chentsov theorem lets us construct a modification of X , i.e. a process X' such that $\forall t. X_t = X'_t$ almost surely, that has Hölder continuous paths.

In this work, we mechanise the Kolmogorov-Chentsov theorem. To get there, we develop a theory of stochastic processes, together with Hölder continuity, convergence in measure, and arbitrary intervals of dyadic rationals.

With this, we pave the way towards a construction of Brownian motion. The work is based on the exposition in Achim Klenke's probability theory text [1].

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1 Supporting lemmas

theory *Kolmogorov-Chentsov-Extras*
imports *HOL-Probability.Probability*
begin

lemma *atLeastAtMost-induct*[*consumes 1, case-names base Suc*]:
assumes $x \in \{n..m\}$
and $P\ n$
and $\bigwedge k. \llbracket k \geq n; k < m; P\ k \rrbracket \implies P\ (Suc\ k)$
shows $P\ x$
 $\langle proof \rangle$

lemma *eventually-prodI'*:
assumes *eventually* $P\ F$ *eventually* $Q\ G\ \forall x\ y. P\ x \longrightarrow Q\ y \longrightarrow R\ (x,y)$
shows *eventually* $R\ (F \times_F G)$
 $\langle proof \rangle$

Analogous to $\llbracket almost\ everywhere\ ?M\ ?P; \bigwedge N. \llbracket \bigwedge x. x \in space\ ?M - N \implies ?P\ x; N \in null\ sets\ ?M \rrbracket \implies ?thesis \rrbracket \implies ?thesis$

lemma *AE-I3*:
assumes $\bigwedge x. x \in space\ M - N \implies P\ x\ N \in null\ sets\ M$
shows *AE* x *in* $M. P\ x$
 $\langle proof \rangle$

Extends $\llbracket (?f \longrightarrow ?l)\ ?F; (?g \longrightarrow ?m)\ ?F \rrbracket \implies ((\lambda x. dist\ (?f\ x)\ (?g\ x)) \longrightarrow dist\ ?l\ ?m)\ ?F$

lemma *tendsto-dist-prod*:
fixes $l\ m :: 'a :: metric\ space$
assumes $f: (f \longrightarrow l)\ F$
and $g: (g \longrightarrow m)\ G$
shows $((\lambda x. dist\ (f\ (fst\ x))\ (g\ (snd\ x))) \longrightarrow dist\ l\ m)\ (F \times_F G)$
 $\langle proof \rangle$

lemma *borel-measurable-at-within-metric* [*measurable*]:
fixes $f :: 'a :: first\ countable\ topology \Rightarrow 'b \Rightarrow 'c :: metric\ space$
assumes [*measurable*]: $\bigwedge t. t \in S \implies f\ t \in borel\ measurable\ M$
and *convergent*: $\bigwedge \omega. \omega \in space\ M \implies \exists l. ((\lambda t. f\ t\ \omega) \longrightarrow l)$ (*at* x *within* S)
and *nontrivial*: *at* x *within* $S \neq \perp$
shows $(\lambda \omega. Lim\ (at\ x\ within\ S)\ (\lambda t. f\ t\ \omega)) \in borel\ measurable\ M$
 $\langle proof \rangle$

lemma *Max-finite-image-ex*:
assumes *finite* $S\ S \neq \{\}$ $P\ (MAX\ k \in S. f\ k)$
shows $\exists k \in S. P\ (f\ k)$
 $\langle proof \rangle$

lemma *measure-leq-emeasure-ennreal*: $0 \leq x \implies emeasure\ M\ A \leq ennreal\ x \implies measure\ M\ A \leq x$

<proof>

lemma *Union-add-subset*: $(m :: \text{nat}) \leq n \implies (\bigcup k. A (k + n)) \subseteq (\bigcup k. A (k + m))$
<proof>

lemma *floor-in-Nats* [simp]: $x \geq 0 \implies \lfloor x \rfloor \in \mathbb{N}$
<proof>

lemma *triangle-ineq-list*:
fixes $l :: ('a :: \text{metric-space}) \text{ list}$
assumes $l \neq []$
shows $\text{dist} (\text{hd } l) (\text{last } l) \leq (\sum_{i=1..length\ l - 1} \text{dist } (l[i]) (l[i+1]))$
<proof>

lemma *triangle-ineq-sum*:
fixes $f :: \text{nat} \Rightarrow 'a :: \text{metric-space}$
assumes $n \leq m$
shows $\text{dist} (f\ n) (f\ m) \leq (\sum_{i=\text{Suc } n..m} \text{dist} (f\ i) (f\ (i-1)))$
<proof>

lemma (in *product-prob-space*) *indep-vars-PiM-coordinate*:
assumes $I \neq \{\}$
shows $\text{prob-space.indep-vars } (\prod_M i \in I. M\ i)\ M (\lambda x\ f. f\ x)\ I$
<proof>

lemma (in *prob-space*) *indep-sets-indep-set*:
assumes $\text{indep-sets } F\ I\ i \in I\ j \in I\ i \neq j$
shows $\text{indep-set } (F\ i)\ (F\ j)$
<proof>

lemma (in *prob-space*) *indep-vars-indep-var*:
assumes $\text{indep-vars } M'\ X\ I\ i \in I\ j \in I\ i \neq j$
shows $\text{indep-var } (M'\ i)\ (X\ i)\ (M'\ j)\ (X\ j)$
<proof>

end

2 Intervals of dyadic rationals

theory *Dyadic-Interval*
imports *HOL-Analysis.Analysis*
begin

In this file we describe intervals of dyadic numbers $S..T$ for reals $S\ T$. We use the floor and ceiling functions to approximate the numbers with increasing accuracy.

lemma *frac-floor*: $\lfloor x \rfloor = x - \text{frac } x$

<proof>

lemma *frac-ceil*: $\lceil x \rceil = x + \text{frac}(-x)$
<proof>

lemma *floor-pow2-lim*: $(\lambda n. \lfloor 2^{\wedge} n * T \rfloor / 2^{\wedge} n) \longrightarrow T$
<proof>

lemma *floor-pow2-leq*: $\lfloor 2^{\wedge} n * T \rfloor / 2^{\wedge} n \leq T$
<proof>

lemma *ceil-pow2-lim*: $(\lambda n. \lceil 2^{\wedge} n * T \rceil / 2^{\wedge} n) \longrightarrow T$
<proof>

lemma *ceil-pow2-geq*: $\lceil 2^{\wedge} n * T \rceil / 2^{\wedge} n \geq T$
<proof>

dyadic_interval_step n S T is the collection of dyadic numbers in $\{S..T\}$ with denominator 2^n . As $n \rightarrow \infty$ this collection approximates $\{S..T\}$. Compare with *dyadics* $\equiv \bigcup_{k \ m} \{of\text{-nat } m / (2::?'a)^k\}$

definition *dyadic-interval-step* :: $nat \Rightarrow real \Rightarrow real \Rightarrow real \text{ set}$
where *dyadic-interval-step* n S $T \equiv (\lambda k. k / (2^{\wedge} n)) ' \{ \lceil 2^{\wedge} n * S \rceil .. \lfloor 2^{\wedge} n * T \rfloor \}$

definition *dyadic-interval* :: $real \Rightarrow real \Rightarrow real \text{ set}$
where *dyadic-interval* S $T \equiv (\bigcup n. \text{dyadic-interval-step } n \ S \ T)$

lemma *dyadic-interval-step-empty[simp]*: $T < S \implies \text{dyadic-interval-step } n \ S \ T = \{\}$
<proof>

lemma *dyadic-interval-step-singleton[simp]*: $X \in \mathbb{Z} \implies \text{dyadic-interval-step } n \ X \ X = \{X\}$
<proof>

lemma *dyadic-interval-step-zero [simp]*: $\text{dyadic-interval-step } 0 \ S \ T = \text{real-of-int } \{ \lceil S \rceil .. \lfloor T \rfloor \}$
<proof>

lemma *dyadic-interval-step-mem [intro]*:
assumes $x \geq 0 \ T \geq 0 \ x \leq T$
shows $\lfloor 2^{\wedge} n * x \rfloor / 2^{\wedge} n \in \text{dyadic-interval-step } n \ 0 \ T$
<proof>

lemma *dyadic-interval-step-iff*:
 $x \in \text{dyadic-interval-step } n \ S \ T \iff$
 $(\exists k. k \geq \lceil 2^{\wedge} n * S \rceil \wedge k \leq \lfloor 2^{\wedge} n * T \rfloor \wedge x = k / 2^{\wedge} n)$
<proof>

lemma *dyadic-interval-step-memI [intro]*:

assumes $\exists k::int. x = k/2^{\wedge n} x \geq S \ x \leq T$
shows $x \in \text{dyadic-interval-step } n \ S \ T$
 <proof>

lemma *mem-dyadic-interval*: $x \in \text{dyadic-interval } S \ T \longleftrightarrow (\exists n. x \in \text{dyadic-interval-step } n \ S \ T)$
 <proof>

lemma *mem-dyadic-intervalI*: $\exists n. x \in \text{dyadic-interval-step } n \ S \ T \Longrightarrow x \in \text{dyadic-interval } S \ T$
 <proof>

lemma *dyadic-step-leq*: $x \in \text{dyadic-interval-step } n \ S \ T \Longrightarrow x \leq T$
 <proof>

lemma *dyadics-leq*: $x \in \text{dyadic-interval } S \ T \Longrightarrow x \leq T$
 <proof>

lemma *dyadic-step-geq*: $x \in \text{dyadic-interval-step } n \ S \ T \Longrightarrow x \geq S$
 <proof>

lemma *dyadics-geq*: $x \in \text{dyadic-interval } S \ T \Longrightarrow x \geq S$
 <proof>

corollary *dyadic-interval-subset-interval* [simp]: $(\text{dyadic-interval } 0 \ T) \subseteq \{0..T\}$
 <proof>

lemma *zero-in-dyadics*: $T \geq 0 \Longrightarrow 0 \in \text{dyadic-interval-step } n \ 0 \ T$
 <proof>

The following theorem is useful for reasoning with `at_within`

lemma *dyadic-interval-converging-sequence*:
assumes $t \in \{0..T\} \ T \neq 0$
shows $\exists s. \forall n. s \ n \in \text{dyadic-interval } 0 \ T - \{t\} \wedge s \longrightarrow t$
 <proof>

lemma *dyadic-interval-dense*: $\text{closure } (\text{dyadic-interval } 0 \ T) = \{0..T\}$
 <proof>

corollary *dyadic-interval-islimgt*:
assumes $T > 0 \ t \in \{0..T\}$
shows $t \text{ islimpt } \text{dyadic-interval } 0 \ T$
 <proof>

corollary *at-within-dyadic-interval-nontrivial*[simp]:
assumes $T > 0 \ t \in \{0..T\}$
shows $(\text{at } t \text{ within } \text{dyadic-interval } 0 \ T) \neq \text{bot}$
 <proof>

lemma *dyadic-interval-step-finite*[simp]: *finite* (*dyadic-interval-step* n S T)
 ⟨*proof*⟩

lemma *dyadic-interval-countable*[simp]: *countable* (*dyadic-interval* S T)
 ⟨*proof*⟩

lemma *floor-pow2-add-leq*:
 fixes $T :: \text{real}$
 shows $\lfloor 2^n * T \rfloor / 2^n \leq \lfloor 2^{n+k} * T \rfloor / 2^{n+k}$
 ⟨*proof*⟩

corollary *floor-pow2-mono*: *mono* ($\lambda n. \lfloor 2^n * (T :: \text{real}) \rfloor / 2^n$)
 ⟨*proof*⟩

lemma *dyadic-interval-step-Max*: $T \geq 0 \implies \text{Max} (\text{dyadic-interval-step } n \ 0 \ T) = \lfloor 2^n * T \rfloor / 2^n$
 ⟨*proof*⟩

lemma *dyadic-interval-step-subset*:
 $n \leq m \implies \text{dyadic-interval-step } n \ 0 \ T \subseteq \text{dyadic-interval-step } m \ 0 \ T$
 ⟨*proof*⟩

corollary *dyadic-interval-step-mono*:
 assumes $x \in \text{dyadic-interval-step } n \ 0 \ T$ $n \leq m$
 shows $x \in \text{dyadic-interval-step } m \ 0 \ T$
 ⟨*proof*⟩

lemma *dyadic-as-natural*:
 assumes $x \in \text{dyadic-interval-step } n \ 0 \ T$
 shows $\exists! k. x = \text{real } k / 2^n$
 ⟨*proof*⟩

lemma *dyadic-of-natural*:
 assumes $\text{real } k / 2^n \leq T$
 shows $\text{real } k / 2^n \in \text{dyadic-interval-step } n \ 0 \ T$
 ⟨*proof*⟩

lemma *dyadic-interval-minus*:
 assumes $x \in \text{dyadic-interval-step } n \ 0 \ T$ $y \in \text{dyadic-interval-step } n \ 0 \ T$ $x \leq y$
 shows $y - x \in \text{dyadic-interval-step } n \ 0 \ T$
 ⟨*proof*⟩

lemma *dyadic-times-nat*: $x \in \text{dyadic-interval-step } n \ 0 \ T \implies (x * 2^n) \in \mathbb{N}$
 ⟨*proof*⟩

definition *dyadic-expansion* $x \ n \ b \ k \equiv \text{set } b \subseteq \{0,1\}$
 $\wedge \text{length } b = n \wedge x = \text{real-of-int } k + (\sum_{m \in \{1..n\}} \text{real } (b ! (m-1)) / 2^m)$

lemma *dyadic-expansionI*:

assumes $set\ b \subseteq \{0,1\}$ $length\ b = n$ $x = k + (\sum_{m \in \{1..n\}} (b ! (m-1)) / 2^{\wedge} m)$
shows *dyadic-expansion* $x\ n\ b\ k$
<proof>

lemma *dyadic-expansionD*:
assumes *dyadic-expansion* $x\ n\ b\ k$
shows $set\ b \subseteq \{0,1\}$
and $length\ b = n$
and $x = k + (\sum_{m \in \{1..n\}} (b ! (m-1)) / 2^{\wedge} m)$
<proof>

lemma *dyadic-expansion-ex*:
assumes $x \in$ *dyadic-interval-step* $n\ 0\ T$
shows $\exists b\ k.$ *dyadic-expansion* $x\ n\ b\ k$
<proof>

lemma *dyadic-expansion-frac-le-1*:
assumes *dyadic-expansion* $x\ n\ b\ k$
shows $(\sum_{m \in \{1..n\}} (b ! (m-1)) / 2^{\wedge} m) < 1$
<proof>

lemma *dyadic-expansion-frac-range*:
assumes *dyadic-expansion* $x\ n\ b\ k\ m \in \{1..n\}$
shows $b ! (m-1) \in \{0,1\}$
<proof>

lemma *dyadic-expansion-interval*:
assumes *dyadic-expansion* $x\ n\ b\ k\ x \in \{S..T\}$
shows $x \in$ *dyadic-interval-step* $n\ S\ T$
<proof>

lemma *dyadic-expansion-nth-geq*:
assumes *dyadic-expansion* $x\ n\ b\ k\ m \in \{1..n\}\ b ! (m-1) = 1$
shows $x \geq k + 1/2^{\wedge} m$
<proof>

lemma *dyadic-expansion-frac-geq-0*:
assumes *dyadic-expansion* $x\ n\ b\ k$
shows $(\sum_{m \in \{1..n\}} (b ! (m-1)) / 2^{\wedge} m) \geq 0$
<proof>

lemma *dyadic-expansion-frac*:
assumes *dyadic-expansion* $x\ n\ b\ k$
shows $frac\ x = (\sum_{m \in \{1..n\}} (b ! (m-1)) / 2^{\wedge} m)$
<proof>

lemma *dyadic-expansion-floor*:
assumes *dyadic-expansion* $x\ n\ b\ k$

shows $k = \lfloor x \rfloor$
 ⟨proof⟩

lemma *sum-interval-pow2-inv*: $(\sum m \in \{Suc\ l..n\}. (1 :: real) / 2^{\wedge} m) = 1 / 2^{\wedge} l - 1 / 2^{\wedge} n$ **if** $l < n$
 ⟨proof⟩

lemma *dyadic-expansion-unique*:
assumes *dyadic-expansion* $x\ n\ b\ k$
and *dyadic-expansion* $x\ n\ c\ j$
shows $b = c \wedge j = k$
 ⟨proof⟩

end

3 Hölder continuity

theory *Holder-Continuous*
imports *HOL-Analysis.Analysis*
begin

Hölder continuity is a weaker version of Lipschitz continuity.

definition *holder-at-within* :: $real \Rightarrow 'a\ set \Rightarrow 'a \Rightarrow ('a :: metric-space \Rightarrow 'b :: metric-space) \Rightarrow bool$ **where**
holder-at-within $\gamma\ D\ r\ \varphi \equiv \gamma \in \{0 <..1\} \wedge$
 $(\exists \varepsilon > 0. \exists C \geq 0. \forall s \in D. dist\ r\ s < \varepsilon \longrightarrow dist\ (\varphi\ r)\ (\varphi\ s) \leq C * dist\ r\ s\ powr\ \gamma)$

definition *local-holder-on* :: $real \Rightarrow 'a :: metric-space\ set \Rightarrow ('a \Rightarrow 'b :: metric-space) \Rightarrow bool$ **where**
local-holder-on $\gamma\ D\ \varphi \equiv \gamma \in \{0 <..1\} \wedge$
 $(\forall t \in D. \exists \varepsilon > 0. \exists C \geq 0. (\forall r \in D. \forall s \in D. dist\ s\ t < \varepsilon \wedge dist\ r\ t < \varepsilon \longrightarrow dist\ (\varphi\ r)\ (\varphi\ s) \leq C * dist\ r\ s\ powr\ \gamma))$

definition *holder-on* :: $real \Rightarrow 'a :: metric-space\ set \Rightarrow ('a \Rightarrow 'b :: metric-space) \Rightarrow bool$ ($--holder'-on\ 1000$) **where**
holder-on $D\ \varphi \longleftrightarrow \gamma \in \{0 <..1\} \wedge (\exists C \geq 0. (\forall r \in D. \forall s \in D. dist\ (\varphi\ r)\ (\varphi\ s) \leq C * dist\ r\ s\ powr\ \gamma))$

lemma *holder-onI*:
assumes $\gamma \in \{0 <..1\} \exists C \geq 0. (\forall r \in D. \forall s \in D. dist\ (\varphi\ r)\ (\varphi\ s) \leq C * dist\ r\ s\ powr\ \gamma)$
shows *holder-on* $D\ \varphi$
 ⟨proof⟩

We prove various equivalent formulations of local holder continuity, using open and closed balls and inequalities.

lemma *local-holder-on-cball*:

local-holder-on $\gamma D \varphi \longleftrightarrow \gamma \in \{0<..1\} \wedge$
 $(\forall t \in D. \exists \varepsilon > 0. \exists C \geq 0. (\forall r \in \text{cball } t \ \varepsilon \cap D. \forall s \in \text{cball } t \ \varepsilon \cap D. \text{dist } (\varphi r) (\varphi s)$
 $\leq C * \text{dist } r s \text{ powr } \gamma))$
(is ?L \longleftrightarrow ?R)
 <proof>

corollary *local-holder-on-leq-def*: *local-holder-on* $\gamma D \varphi \longleftrightarrow \gamma \in \{0<..1\} \wedge$
 $(\forall t \in D. \exists \varepsilon > 0. \exists C \geq 0. (\forall r \in D. \forall s \in D. \text{dist } s t \leq \varepsilon \wedge \text{dist } r t \leq \varepsilon \longrightarrow \text{dist } (\varphi$
 $r) (\varphi s) \leq C * \text{dist } r s \text{ powr } \gamma))$
 <proof>

corollary *local-holder-on-ball*: *local-holder-on* $\gamma D \varphi \longleftrightarrow \gamma \in \{0<..1\} \wedge$
 $(\forall t \in D. \exists \varepsilon > 0. \exists C \geq 0. (\forall r \in \text{ball } t \ \varepsilon \cap D. \forall s \in \text{ball } t \ \varepsilon \cap D. \text{dist } (\varphi r) (\varphi s)$
 $\leq C * \text{dist } r s \text{ powr } \gamma))$
 <proof>

lemma *local-holder-on-altdef*:
assumes $D \neq \{\}$
shows *local-holder-on* $\gamma D \varphi = (\forall t \in D. (\exists \varepsilon > 0. (\gamma\text{-holder-on } ((\text{cball } t \ \varepsilon) \cap$
 $D) \varphi)))$
 <proof>

lemma *local-holder-on-cong*[cong]:
assumes $\gamma = \varepsilon C = D \wedge x. x \in C \implies \varphi x = \psi x$
shows *local-holder-on* $\gamma C \varphi \longleftrightarrow \text{local-holder-on } \varepsilon D \psi$
 <proof>

lemma *local-holder-onI*:
assumes $\gamma \in \{0<..1\} (\forall t \in D. \exists \varepsilon > 0. \exists C \geq 0. (\forall r \in D. \forall s \in D. \text{dist } s t < \varepsilon \wedge$
 $\text{dist } r t < \varepsilon \longrightarrow \text{dist } (\varphi r) (\varphi s) \leq C * \text{dist } r s \text{ powr } \gamma))$
shows *local-holder-on* $\gamma D \varphi$
 <proof>

lemma *local-holder-ballI*:
assumes $\gamma \in \{0<..1\}$
and $\bigwedge t. t \in D \implies \exists \varepsilon > 0. \exists C \geq 0. \forall r \in \text{ball } t \ \varepsilon \cap D. \forall s \in \text{ball } t \ \varepsilon \cap D.$
 $\text{dist } (\varphi r) (\varphi s) \leq C * \text{dist } r s \text{ powr } \gamma$
shows *local-holder-on* $\gamma D \varphi$
 <proof>

lemma *local-holder-onE*:
assumes *local-holder*: *local-holder-on* $\gamma D \varphi$
and gamma: $\gamma \in \{0<..1\}$
and $t \in D$
obtains εC **where** $\varepsilon > 0 C \geq 0$
 $\bigwedge r s. r \in \text{ball } t \ \varepsilon \cap D \implies s \in \text{ball } t \ \varepsilon \cap D \implies \text{dist } (\varphi r) (\varphi s) \leq C * \text{dist } r$
 $s \text{ powr } \gamma$
 <proof>

Holder continuity matches up with the existing definitions in *HOL-Analysis.Lipschitz*

lemma *holder-1-eq-lipschitz*: 1 -holder-on D $\varphi = (\exists C. \text{lipschitz-on } C D \varphi)$
<proof>

lemma *local-holder-1-eq-local-lipschitz*:
assumes $T \neq \{\}$
shows *local-holder-on* $1 D \varphi = \text{local-lipschitz } T D (\lambda-. \varphi)$
<proof>

lemma *local-holder-refine*:
assumes g : *local-holder-on* $g D \varphi$ $g \leq 1$
and h : $h \leq g$ $h > 0$
shows *local-holder-on* $h D \varphi$
<proof>

lemma *holder-uniform-continuous*:
assumes γ -holder-on $X \varphi$
shows *uniformly-continuous-on* $X \varphi$
<proof>

corollary *holder-on-continuous-on*: γ -holder-on $X \varphi \implies \text{continuous-on } X \varphi$
<proof>

lemma *holder-implies-local-holder*: γ -holder-on $D \varphi \implies \text{local-holder-on } \gamma D \varphi$
<proof>

lemma *local-holder-imp-continuous*:
assumes *local-holder*: *local-holder-on* $\gamma X \varphi$
shows *continuous-on* $X \varphi$
<proof>

lemma *local-holder-compact-imp-holder*:
assumes *compact* I *local-holder-on* $\gamma I \varphi$
shows γ -holder-on $I \varphi$
<proof>

lemma *holder-const*: γ -holder-on $C (\lambda-. c) \longleftrightarrow \gamma \in \{0 <..1\}$
<proof>

lemma *local-holder-const*: *local-holder-on* $\gamma C (\lambda-. c) \longleftrightarrow \gamma \in \{0 <..1\}$
<proof>

end

4 Convergence in measure

theory *Measure-Convergence*
imports *HOL-Probability.Probability*
begin

We use `measure` rather than `emeasure` because `ennreal` is not a metric space, which we need to reason about convergence. By intersecting with the set of finite measure `A`, we don't run into issues where infinity is collapsed to 0. For finite measures this definition is equal to the definition without set `A` – see below.

definition `tendsto-measure` :: `'b measure` \Rightarrow (`'a` \Rightarrow `'b` \Rightarrow (`'c` :: $\{\text{second-countable-topology, metric-space}\}$))
 \Rightarrow (`'b` \Rightarrow `'c`) \Rightarrow `'a filter` \Rightarrow `bool`
where `tendsto-measure` `M X l F` \equiv ($\forall n. X\ n \in \text{borel-measurable } M$) \wedge `l` \in `borel-measurable` `M` \wedge
 $(\forall \varepsilon > 0. \forall A \in \text{fmeasurable } M.$
 $((\lambda n. \text{measure } M (\{\omega \in \text{space } M. \text{dist } (X\ n\ \omega) (l\ \omega) > \varepsilon\} \cap A)) \longrightarrow 0) F)$

abbreviation (in `prob-space`) `tendsto-prob` (**infixr** \longrightarrow_P 55) **where**
 $(f \longrightarrow_P l) F \equiv \text{tendsto-measure } M\ f\ l\ F$

lemma `tendsto-measure-measurable[measurable-dest]`:
`tendsto-measure` `M X l F` \Longrightarrow `X n` \in `borel-measurable` `M`
 $\langle \text{proof} \rangle$

lemma `tendsto-measure-measurable-lim[measurable-dest]`:
`tendsto-measure` `M X l F` \Longrightarrow `l` \in `borel-measurable` `M`
 $\langle \text{proof} \rangle$

lemma `tendsto-measure-mono`: `F` \leq `F'` \Longrightarrow `tendsto-measure` `M f l F'` \Longrightarrow `tendsto-measure` `M f l F`
 $\langle \text{proof} \rangle$

lemma `tendsto-measureI`:
assumes `[measurable]`: $\bigwedge n. X\ n \in \text{borel-measurable } M$ `l` \in `borel-measurable` `M`
and $\bigwedge \varepsilon A. \varepsilon > 0 \Longrightarrow A \in \text{fmeasurable } M \Longrightarrow$
 $((\lambda n. \text{measure } M (\{\omega \in \text{space } M. \text{dist } (X\ n\ \omega) (l\ \omega) > \varepsilon\} \cap A)) \longrightarrow 0) F$
shows `tendsto-measure` `M X l F`
 $\langle \text{proof} \rangle$

lemma (in `finite-measure`) `finite-tendsto-measureI`:
assumes `[measurable]`: $\bigwedge n. f'\ n \in \text{borel-measurable } M$ `f` \in `borel-measurable` `M`
and $\bigwedge \varepsilon. \varepsilon > 0 \Longrightarrow ((\lambda n. \text{measure } M \{\omega \in \text{space } M. \text{dist } (f'\ n\ \omega) (f\ \omega) > \varepsilon\}) \longrightarrow 0) F$
shows `tendsto-measure` `M f' f F`
 $\langle \text{proof} \rangle$

lemma (in `finite-measure`) `finite-tendsto-measureD`:
assumes `[measurable]`: `tendsto-measure` `M f' f F`
shows ($\forall \varepsilon > 0. ((\lambda n. \text{measure } M \{\omega \in \text{space } M. \text{dist } (f'\ n\ \omega) (f\ \omega) > \varepsilon\}) \longrightarrow 0) F$)
 $\langle \text{proof} \rangle$

lemma (in `finite-measure`) `tendsto-measure-leq`:

assumes $[measurable]: \bigwedge n. f' n \in \text{borel-measurable } M \ f \in \text{borel-measurable } M$
shows $\text{tendsto-measure } M \ f' \ f \ F \longleftrightarrow$
 $(\forall \varepsilon > 0. ((\lambda n. \text{measure } M \ \{\omega \in \text{space } M. \text{dist } (f' n \ \omega) \ (f \ \omega) \geq \varepsilon\}) \longrightarrow 0)$
 $F) \text{ (is } ?L \longleftrightarrow ?R)$
 $\langle \text{proof} \rangle$

abbreviation $\text{LIMSEQ-measure } M \ f \ l \equiv \text{tendsto-measure } M \ f \ l \ \text{sequentially}$

lemma $\text{LIMSEQ-measure-def}: \text{LIMSEQ-measure } M \ f \ l \longleftrightarrow$
 $(\forall n. f \ n \in \text{borel-measurable } M) \wedge (l \in \text{borel-measurable } M) \wedge$
 $(\forall \varepsilon > 0. \forall A \in \text{fmeasurable } M.$
 $(\lambda n. \text{measure } M \ (\{\omega \in \text{space } M. \text{dist } (f \ n \ \omega) \ (l \ \omega) > \varepsilon\} \cap A)) \longrightarrow 0)$
 $\langle \text{proof} \rangle$

lemma LIMSEQ-measureD :
assumes $\text{LIMSEQ-measure } M \ f \ l \ \varepsilon > 0 \ A \in \text{fmeasurable } M$
shows $(\lambda n. \text{measure } M \ (\{\omega \in \text{space } M. \text{dist } (f \ n \ \omega) \ (l \ \omega) > \varepsilon\} \cap A)) \longrightarrow 0$
 $\langle \text{proof} \rangle$

lemma fmeasurable-inter : $\llbracket A \in \text{sets } M; B \in \text{fmeasurable } M \rrbracket \implies A \cap B \in \text{fmeasurable } M$
 $\langle \text{proof} \rangle$

lemma $\text{LIMSEQ-measure-emeasure}$:
assumes $\text{LIMSEQ-measure } M \ f \ l \ \varepsilon > 0 \ A \in \text{fmeasurable } M$
and $[measurable]: \bigwedge i. f \ i \in \text{borel-measurable } M \ l \in \text{borel-measurable } M$
shows $(\lambda n. \text{emeasure } M \ (\{\omega \in \text{space } M. \text{dist } (f \ n \ \omega) \ (l \ \omega) > \varepsilon\} \cap A)) \longrightarrow 0$
 $\langle \text{proof} \rangle$

lemma $\text{measure-Lim-within-LIMSEQ}$:
fixes $a :: 'a :: \text{first-countable-topology}$
assumes $\bigwedge t. X \ t \in \text{borel-measurable } M \ L \in \text{borel-measurable } M$
assumes $\bigwedge S. \llbracket (\forall n. S \ n \neq a \wedge S \ n \in T); S \longrightarrow a \rrbracket \implies \text{LIMSEQ-measure } M$
 $(\lambda n. X \ (S \ n)) \ L$
shows $\text{tendsto-measure } M \ X \ L \ \text{(at } a \ \text{within } T)$
 $\langle \text{proof} \rangle$

definition $\text{tendsto-AE} :: 'b \ \text{measure} \Rightarrow ('a \Rightarrow 'b \Rightarrow 'c :: \text{topological-space}) \Rightarrow ('b$
 $\Rightarrow 'c) \Rightarrow 'a \ \text{filter} \Rightarrow \text{bool}$ **where**
 $\text{tendsto-AE } M \ f' \ l \ F \longleftrightarrow (\text{AE } \omega \ \text{in } M. ((\lambda n. f' \ n \ \omega) \longrightarrow l \ \omega) \ F)$

lemma $\text{LIMSEQ-ae-pointwise}$: $(\bigwedge x. (\lambda n. f \ n \ x) \longrightarrow l \ x) \implies \text{tendsto-AE } M \ f \ l$
 sequentially
 $\langle \text{proof} \rangle$

lemma $\text{tendsto-AE-within-LIMSEQ}$:
fixes $a :: 'a :: \text{first-countable-topology}$
assumes $\bigwedge S. \llbracket (\forall n. S \ n \neq a \wedge S \ n \in T); S \longrightarrow a \rrbracket \implies \text{tendsto-AE } M \ (\lambda n.$
 $X \ (S \ n)) \ L \ \text{sequentially}$

shows $\text{tendsto-AE } M X L$ (at a within T)
 ⟨proof⟩

lemma *LIMSEQ-dominated-convergence*:
fixes $X :: \text{nat} \Rightarrow \text{real}$
assumes $X \longrightarrow L$ ($\bigwedge n. Y n \leq X n$) ($\bigwedge n. Y n \geq L$)
shows $Y \longrightarrow L$
 ⟨proof⟩

Klenke remark 6.4

lemma *measure-conv-imp-AE-sequentially*:
assumes [*measurable*]: $\bigwedge n. f' n \in \text{borel-measurable } M$ $f \in \text{borel-measurable } M$
and $\text{tendsto-AE } M f' f$ *sequentially*
shows *LIMSEQ-measure* $M f' f$
 ⟨proof⟩

corollary *LIMSEQ-measure-pointwise*:
assumes $\bigwedge x. (\lambda n. f n x) \longrightarrow f' x$ $\bigwedge n. f n \in \text{borel-measurable } M$ $f' \in \text{borel-measurable } M$
shows *LIMSEQ-measure* $M f f'$
 ⟨proof⟩

lemma *Lim-measure-pointwise*:
fixes $a :: 'a :: \text{first-countable-topology}$
assumes $\bigwedge x. ((\lambda n. f n x) \longrightarrow f' x)$ (at a within T) $\bigwedge n. f n \in \text{borel-measurable } M$ $f' \in \text{borel-measurable } M$
shows $\text{tendsto-measure } M f f'$ (at a within T)
 ⟨proof⟩

corollary *measure-conv-imp-AE-at-within*:
fixes $x :: 'a :: \text{first-countable-topology}$
assumes [*measurable*]: $\bigwedge n. f' n \in \text{borel-measurable } M$ $f \in \text{borel-measurable } M$
and $\text{tendsto-AE } M f' f$ (at x within S)
shows $\text{tendsto-measure } M f' f$ (at x within S)
 ⟨proof⟩

Klenke remark 6.5

lemma (in *sigma-finite-measure*) *LIMSEQ-measure-unique-AE*:
fixes $f :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{second-countable-topology, metric-space}\}$
assumes [*measurable*]: $\bigwedge n. f n \in \text{borel-measurable } M$ $l \in \text{borel-measurable } M$ $l' \in \text{borel-measurable } M$
and *LIMSEQ-measure* $M f l$ *LIMSEQ-measure* $M f l'$
shows *AE* x in $M. l x = l' x$
 ⟨proof⟩

corollary (in *sigma-finite-measure*) *LIMSEQ-ae-unique-AE*:
fixes $f :: \text{nat} \Rightarrow 'a \Rightarrow 'b :: \{\text{second-countable-topology, metric-space}\}$
assumes $\bigwedge n. f n \in \text{borel-measurable } M$ $l \in \text{borel-measurable } M$ $l' \in \text{borel-measurable } M$
shows *AE* x in $M. l x = l' x$

and *tendsto-AE M f l sequentially tendsto-AE M f l' sequentially*
shows *AE x in M. l x = l' x*
 ⟨*proof*⟩

lemma (in *sigma-finite-measure*) *tendsto-measure-at-within-eq-AE*:
fixes *f :: 'b :: first-countable-topology ⇒ 'a ⇒ 'c :: {second-countable-topology,metric-space}*
assumes *f-measurable: $\bigwedge x. x \in S \implies f x \in \text{borel-measurable } M$*
and *l-measurable: $l \in \text{borel-measurable } M$ $l' \in \text{borel-measurable } M$*
and *tendsto: tendsto-measure M f l (at t within S) tendsto-measure M f l' (at t within S)*
and *not-bot: (at t within S) $\neq \perp$*
shows *AE x in M. l x = l' x*
 ⟨*proof*⟩
end

5 Stochastic processes

theory *Stochastic-Processes*
imports *Kolmogorov-Chentsov-Extras Dyadic-Interval*
begin

A stochastic process is an indexed collection of random variables. For compatibility with `product_prob_space` we don't enforce conditions on the index set I in the assumptions.

locale *stochastic-process = prob-space +*
fixes *M' :: 'b measure*
and *I :: 't set*
and *X :: 't ⇒ 'a ⇒ 'b*
assumes *random-process[measurable]: $\bigwedge i. \text{random-variable } M' (X i)$*

sublocale *stochastic-process \subseteq product: product-prob-space ($\lambda t. \text{distr } M M' (X t)$)*
 ⟨*proof*⟩

lemma (in *prob-space*) *stochastic-process I*:
assumes *$\bigwedge i. \text{random-variable } M' (X i)$*
shows *stochastic-process M M' X*
 ⟨*proof*⟩

typedef (*'t, 'a, 'b*) *stochastic-process =*
 {*(M :: 'a measure, M' :: 'b measure, I :: 't set, X :: 't ⇒ 'a ⇒ 'b).*
stochastic-process M M' X}
 ⟨*proof*⟩

setup-lifting *type-definition-stochastic-process*

lift-definition *proc-source :: ('t,'a,'b) stochastic-process ⇒ 'a measure*
is *fst* ⟨*proof*⟩

interpretation *proc-source*: *prob-space proc-source X*
 ⟨*proof*⟩

lift-definition *proc-target* :: ('t,'a,'b) *stochastic-process* ⇒ 'b *measure*
 is *fst* ∘ *snd* ⟨*proof*⟩

lift-definition *proc-index* :: ('t,'a,'b) *stochastic-process* ⇒ 't *set*
 is *fst* ∘ *snd* ∘ *snd* ⟨*proof*⟩

lift-definition *process* :: ('t,'a,'b) *stochastic-process* ⇒ 't ⇒ 'a ⇒ 'b
 is *snd* ∘ *snd* ∘ *snd* ⟨*proof*⟩

declare [[*coercion process*]]

lemma *stochastic-process-construct* [*simp*]: *stochastic-process (proc-source X) (proc-target X) (process X)*
 ⟨*proof*⟩

interpretation *stochastic-process proc-source X proc-target X proc-index X process X*
 ⟨*proof*⟩

lemma *stochastic-process-measurable* [*measurable*]: *process X t ∈ (proc-source X) →_M (proc-target X)*
 ⟨*proof*⟩

Here we construct a process on a given index set. For this we need to produce measurable functions for indices outside the index set; we use the constant function, but it needs to point at an element of the target set to be measurable.

context *prob-space*
begin

lift-definition *process-of* :: 'b *measure* ⇒ 't *set* ⇒ ('t ⇒ 'a ⇒ 'b) ⇒ 'b ⇒ ('t,'a,'b) *stochastic-process*
 is λ *M' I X ω*. if (∀ *t* ∈ *I*. *X t* ∈ *M* →_M *M'*) ∧ *ω* ∈ *space M'*
 then (*M*, *M'*, *I*, (λ*t*. if *t* ∈ *I* then *X t* else (λ-. *ω*)))
 else (return (*sigma UNIV* {}, *UNIV*)) (*SOME x. True*), *sigma UNIV UNIV*,
I, λ- -. *ω*)
 ⟨*proof*⟩

lemma *index-process-of* [*simp*]: *proc-index (process-of M' I X ω) = I*
 ⟨*proof*⟩

lemma

assumes ∀ *t* ∈ *I*. *X t* ∈ *M* →_M *M'* *ω* ∈ *space M'*

shows

source-process-of [*simp*]: *proc-source (process-of M' I X ω) = M* **and**
target-process-of [*simp*]: *proc-target (process-of M' I X ω) = M'* **and**

process-process-of[simp]: process (process-of $M' I X \omega$) = (λt . if $t \in I$ then $X t$ else ($\lambda \cdot$. ω))
<proof>

lemma process-of-apply:

assumes $\forall t \in I. X t \in M \rightarrow_M M' \omega \in \text{space } M' t \in I$

shows *process (process-of $M' I X \omega$) $t = X t$*

<proof>

end

We define the finite-dimensional distributions of our process.

lift-definition distributions :: ('t, 'a, 'b) stochastic-process \Rightarrow 't set \Rightarrow ('t \Rightarrow 'b) measure

is $\lambda(M, M', -, X) T. (\Pi_M t \in T. \text{distr } M M' (X t))$ *<proof>*

lemma distributions-altdef: distributions $X T = (\Pi_M t \in T. \text{distr (proc-source } X) (\text{proc-target } X) (X t))$

<proof>

lemma prob-space-distributions: prob-space (distributions $X J$)

<proof>

lemma sets-distributions: sets (distributions $X J$) = sets (PiM J ($\lambda \cdot$. (proc-target X)))

<proof>

lemma space-distributions: space (distributions $X J$) = ($\Pi_E i \in J. \text{space (proc-target } X)$)

<proof>

lemma emeasure-distributions:

assumes *finite $J \wedge j. j \in J \implies A j \in \text{sets (proc-target } X)$*

shows *emeasure (distributions $X J$) (PiE J A) = ($\prod j \in J. \text{emeasure (distr (proc-source } X) (\text{proc-target } X) (X j)) (A j)$)*

<proof>

interpretation projective-family (proc-index X) distributions $X (\lambda \cdot$. proc-target $X)$

<proof>

locale polish-stochastic = stochastic-process $M \text{ borel} :: 'b::\text{polish-space measure } I X$ for M and I and X

lemma distributed-cong-random-variable:

assumes $M = K N = L A E x \text{ in } M. X x = Y x X \in M \rightarrow_M N Y \in K \rightarrow_M L$
 $f \in \text{borel-measurable } N$

shows *distributed $M N X f \longleftrightarrow$ distributed $K L Y f$*

<proof>

For all sorted lists of indices, the increments specified by this list are independent

lift-definition *indep-increments* :: ('t :: linorder, 'a, 'b :: minus) stochastic-process \Rightarrow bool **is**

$\lambda(M, M', I, X).$
 $(\forall l. \text{set } l \subseteq I \wedge \text{sorted } l \wedge \text{length } l \geq 2 \longrightarrow$
 $\text{prob-space.indep-vars } M (\lambda-. M') (\lambda k v. X (!k) v - X (!k-1)) v) \{1..<\text{length } l\}$
<proof>

lemma *indep-incrementsE*:

assumes *indep-increments* X
and $\text{set } l \subseteq \text{proc-index } X \wedge \text{sorted } l \wedge \text{length } l \geq 2$
shows $\text{prob-space.indep-vars } (\text{proc-source } X) (\lambda-. \text{proc-target } X)$
 $(\lambda k v. X (!k) v - X (!k-1)) v) \{1..<\text{length } l\}$
<proof>

lemma *indep-incrementsI*:

assumes $\bigwedge l. \text{set } l \subseteq \text{proc-index } X \Longrightarrow \text{sorted } l \Longrightarrow \text{length } l \geq 2 \Longrightarrow$
 $\text{prob-space.indep-vars } (\text{proc-source } X) (\lambda-. \text{proc-target } X) (\lambda k v. X (!k) v - X$
 $(!(k-1)) v) \{1..<\text{length } l\}$
shows *indep-increments* X
<proof>

lemma *indep-increments-indep-var*:

assumes *indep-increments* X $h \in \text{proc-index } X$ $j \in \text{proc-index } X$ $k \in \text{proc-index } X$
 $h \leq j \leq k$
shows $\text{prob-space.indep-var } (\text{proc-source } X) (\text{proc-target } X) (\lambda v. X j v - X h v)$
 $(\text{proc-target } X) (\lambda v. X k v - X j v)$
<proof>

definition *stationary-increments* X $\longleftrightarrow (\forall t1 t2 k. t1 > 0 \wedge t2 > 0 \wedge k > 0 \longrightarrow$

$\text{distr } (\text{proc-source } X) (\text{proc-target } X) (\lambda v. X (t1 + k) v - X t1 v) =$
 $\text{distr } (\text{proc-source } X) (\text{proc-target } X) (\lambda v. X (t2 + k) v - X t2 v)$

Processes on the same source measure space, with the same index space, but not necessarily the same target measure since we only care about the measurable target space, not the measure

lift-definition *compatible* :: ('t,'a,'b) stochastic-process \Rightarrow ('t,'a,'b) stochastic-process \Rightarrow bool

is $\lambda(Mx, M'x, Ix, X) (My, M'y, Iy, -). Mx = My \wedge \text{sets } M'x = \text{sets } M'y \wedge Ix = Iy$ *<proof>*

lemma *compatibleI*:

assumes $\text{proc-source } X = \text{proc-source } Y$ $\text{sets } (\text{proc-target } X) = \text{sets } (\text{proc-target } Y)$
 $\text{proc-index } X = \text{proc-index } Y$

shows *compatible* $X Y$
<proof>

lemma

assumes *compatible* $X Y$

shows

compatible-source [*dest*]: *proc-source* $X = \text{proc-source } Y$ **and**
compatible-target [*dest*]: *sets* (*proc-target* X) = *sets* (*proc-target* Y) **and**
compatible-index [*dest*]: *proc-index* $X = \text{proc-index } Y$

<proof>

lemma *compatible-refl* [*simp*]: *compatible* $X X$

<proof>

lemma *compatible-sym*: *compatible* $X Y \implies \text{compatible } Y X$

<proof>

lemma *compatible-trans*:

assumes *compatible* $X Y$ *compatible* $Y Z$

shows *compatible* $X Z$

<proof>

lemma (**in** *prob-space*) *compatible-process-of*:

assumes *measurable*: $\forall t \in I. X t \in M \rightarrow_M M' \forall t \in I. Y t \in M \rightarrow_M M'$
and $a \in \text{space } M' b \in \text{space } M'$

shows *compatible* (*process-of* $M' I X a$) (*process-of* $M' I Y b$)

<proof>

definition *modification* :: $(t, 'a, 'b)$ *stochastic-process* $\Rightarrow (t, 'a, 'b)$ *stochastic-process*
 \Rightarrow *bool* **where**

modification $X Y \iff \text{compatible } X Y \wedge (\forall t \in \text{proc-index } X. AE x \text{ in } \text{proc-source } X. X t x = Y t x)$

lemma *modificationI* [*intro*]:

assumes *compatible* $X Y \wedge t. t \in \text{proc-index } X \implies AE x \text{ in } \text{proc-source } X. X t x = Y t x$

shows *modification* $X Y$

<proof>

lemma *modificationD* [*dest*]:

assumes *modification* $X Y$

shows *compatible* $X Y$

and $\wedge t. t \in \text{proc-index } X \implies AE x \text{ in } \text{proc-source } X. X t x = Y t x$

<proof>

lemma *modification-null-set*:

assumes *modification* $X Y t \in \text{proc-index } X$

obtains N **where** $\{x \in \text{space } (\text{proc-source } X). X t x \neq Y t x\} \subseteq N N \in \text{null-sets } (\text{proc-source } X)$

<proof>

lemma *modification-refl* [*simp*]: *modification X X*
<proof>

lemma *modification-sym*: *modification X Y \implies modification Y X*
<proof>

lemma *modification-trans*:
assumes *modification X Y modification Y Z*
shows *modification X Z*
<proof>

lemma *modification-imp-identical-distributions*:
assumes *modification: modification X Y*
and *index: T \subseteq proc-index X*
shows *distributions X T = distributions Y T*
<proof>

definition *indistinguishable :: ('t,'a,'b) stochastic-process \implies ('t,'a,'b) stochastic-process*
 \implies bool where
indistinguishable X Y \longleftrightarrow compatible X Y \wedge
*($\exists N \in$ null-sets (proc-source X). $\forall t \in$ proc-index X. $\{x \in$ space (proc-source X).
 $X t x \neq Y t x\} \subseteq N$)*

lemma *indistinguishableI*:
assumes *compatible X Y*
and *$\exists N \in$ null-sets (proc-source X). ($\forall t \in$ proc-index X. $\{x \in$ space (proc-source
X). $X t x \neq Y t x\} \subseteq N$)*
shows *indistinguishable X Y*
<proof>

lemma *indistinguishable-null-set*:
assumes *indistinguishable X Y*
obtains *N where*
N \in null-sets (proc-source X)
 $\wedge t. t \in$ proc-index X \implies $\{x \in$ space (proc-source X). $X t x \neq Y t x\} \subseteq N$
<proof>

lemma *indistinguishableD*:
assumes *indistinguishable X Y*
shows *compatible X Y*
and *$\exists N \in$ null-sets (proc-source X). ($\forall t \in$ proc-index X. $\{x \in$ space (proc-source
X). $X t x \neq Y t x\} \subseteq N$)*
<proof>

lemma *indistinguishable-eq-AE*:
assumes *indistinguishable X Y*
shows *AE x in proc-source X. $\forall t \in$ proc-index X. $X t x = Y t x$*

<proof>

lemma *indistinguishable-null-ex:*

assumes *indistinguishable X Y*

shows $\exists N \in \text{null-sets}(\text{proc-source } X). \{x \in \text{space}(\text{proc-source } X). \exists t \in \text{proc-index } X. X t x \neq Y t x\} \subseteq N$

<proof>

lemma *indistinguishable-refl [simp]: indistinguishable X X*

<proof>

lemma *indistinguishable-sym: indistinguishable X Y \implies indistinguishable Y X*

<proof>

lemma *indistinguishable-trans:*

assumes *indistinguishable X Y indistinguishable Y Z*

shows *indistinguishable X Z*

<proof>

lemma *indistinguishable-modification: indistinguishable X Y \implies modification X Y*

<proof>

Klenke 21.5(i)

lemma *modification-countable:*

assumes *modification X Y countable (proc-index X)*

shows *indistinguishable X Y*

<proof>

Klenke 21.5(ii). The textbook statement is more general - we reduce right continuity to regular continuity

lemma *modification-continuous-indistinguishable:*

fixes $X :: (\text{real}, 'a, 'b :: \text{metric-space}) \text{stochastic-process}$

assumes *modification: modification X Y*

and *interval: $\exists T > 0. \text{proc-index } X = \{0..T\}$*

and *rc: AE ω in proc-source X. continuous-on (proc-index X) ($\lambda t. X t \omega$)*

(is AE ω in proc-source X. ?cont-X ω)

AE ω in proc-source Y. continuous-on (proc-index Y) ($\lambda t. Y t \omega$)

(is AE ω in proc-source Y. ?cont-Y ω)

shows *indistinguishable X Y*

<proof>

lift-definition *restrict-index :: ('t, 'a, 'b) stochastic-process \implies 't set \implies ('t, 'a, 'b) stochastic-process*

is $\lambda(M, M', I, X) T. (M, M', T, X)$ *<proof>*

lemma

shows

$restrict-index-source[simp]: proc-source (restrict-index X T) = proc-source X$
and
 $restrict-index-target[simp]: proc-target (restrict-index X T) = proc-target X$ **and**
 $restrict-index-index[simp]: proc-index (restrict-index X T) = T$ **and**
 $restrict-index-process[simp]: process (restrict-index X T) = process X$
 ⟨proof⟩

lemma $restrict-index-override[simp]: restrict-index (restrict-index X T) S = restrict-index X S$
 ⟨proof⟩

lemma $compatible-restrict-index:$
assumes $compatible X Y$
shows $compatible (restrict-index X S) (restrict-index Y S)$
 ⟨proof⟩

lemma $modification-restrict-index:$
assumes $modification X Y S \subseteq proc-index X$
shows $modification (restrict-index X S) (restrict-index Y S)$
 ⟨proof⟩

lemma $indistinguishable-restrict-index:$
assumes $indistinguishable X Y S \subseteq proc-index X$
shows $indistinguishable (restrict-index X S) (restrict-index Y S)$
 ⟨proof⟩

lemma $AE-eq-minus [intro]:$
fixes $a :: 'a \Rightarrow ('b :: real-normed-vector)$
assumes $AE x in M. a x = b x$ $AE x in M. c x = d x$
shows $AE x in M. a x - c x = b x - d x$
 ⟨proof⟩

lemma $modification-indep-increments:$
fixes $X Y :: ('a :: linorder, 'b, 'c :: \{second-countable-topology, real-normed-vector\})$
 $stochastic-process$
assumes $modification X Y sets (proc-target Y) = sets borel$
shows $indep-increments X \implies indep-increments Y$
 ⟨proof⟩

end

6 The Kolmogorov-Chentsov theorem

theory $Kolmogorov-Chentsov$
imports $Stochastic-Processes Holder-Continuous Dyadic-Interval Measure-Convergence$
begin

6.1 Supporting lemmas

The main contribution of this file is the Kolmogorov-Chentsov theorem: given a stochastic process that satisfies some continuity properties, we can construct a Hölder continuous modification. We first prove some auxiliary lemmas before moving on to the main construction.

Klenke 5.11: Markov inequality. Compare with $\llbracket (\lambda x. ?u x * \text{indicator } ?A x) \in \text{borel-measurable } ?M; ?A \in \text{sets } ?M \rrbracket \implies \text{emeasure } ?M \{x \in ?A. 1 \leq ?c * ?u x\} \leq ?c * \text{set-nn-integral } ?M ?A ?u$

lemma *nn-integral-Markov-inequality-extended*:

fixes $f :: \text{real} \Rightarrow \text{ennreal}$ **and** $\varepsilon :: \text{real}$ **and** $X :: 'a \Rightarrow \text{real}$
assumes *mono*: *mono-on* ($\text{range } X \cup \{0<..\}$) f
and *finite*: $\bigwedge x. f x < \infty$
and $e: \varepsilon > 0 f \varepsilon > 0$
and [*measurable*]: $X \in \text{borel-measurable } M$
shows $\text{emeasure } M \{p \in \text{space } M. (X p) \geq \varepsilon\} \leq (\int^+ x. f (X x) \partial M) / f \varepsilon$
<proof>

lemma *nn-integral-Markov-inequality-extended-rnv*:

fixes $f :: \text{real} \Rightarrow \text{real}$ **and** $\varepsilon :: \text{real}$ **and** $X :: 'a \Rightarrow 'b :: \text{real-normed-vector}$
assumes [*measurable*]: $X \in \text{borel-measurable } M$
and *mono*: *mono-on* $\{0<..\}$ f
and $e: \varepsilon > 0 f \varepsilon > 0$
shows $\text{emeasure } M \{p \in \text{space } M. \text{norm } (X p) \geq \varepsilon\} \leq (\int^+ x. f (\text{norm } (X x)) \partial M) / f \varepsilon$
<proof>

6.2 Kolmogorov-Chentsov

Klenke theorem 21.6 - Kolmogorov-Chentsov

locale *Kolmogorov-Chentsov* =

fixes $X :: (\text{real}, 'a, 'b :: \text{polish-space}) \text{stochastic-process}$
and $a b C \gamma :: \text{real}$
assumes *index[simp]*: *proc-index* $X = \{0<..\}$
and *target-borel[simp]*: *proc-target* $X = \text{borel}$
and *gt-0*: $a > 0 b > 0 C > 0$
and *b-leq-a*: $b \leq a$
and *gamma*: $\gamma \in \{0<..\leq b/a\}$
and *expectation*: $\bigwedge s t. \llbracket s \geq 0; t \geq 0 \rrbracket \implies$
 $(\int^+ x. \text{dist } (X t x) (X s x) \text{powr } a \partial \text{proc-source } X) \leq C * \text{dist } t s \text{powr } (1+b)$
begin

lemma *gamma-0-1[simp]*: $\gamma \in \{0<..\leq 1\}$
<proof>

lemma *gamma-gt-0[simp]*: $\gamma > 0$

<proof>

lemma *gamma-le-1[simp]*: $\gamma \leq 1$
<proof>

abbreviation *source* \equiv *proc-source* *X*

lemma *X-borel-measurable[measurable]*: $X \in$ *borel-measurable source* **for** *t*
<proof>

lemma *markov*: $\mathcal{P}(x \text{ in } \textit{source}. \varepsilon \leq \textit{dist} (X \ t \ x) (X \ s \ x)) \leq (C * \textit{dist} \ t \ s \ \textit{powr} (1 + b)) / \varepsilon \ \textit{powr} \ a$
if $s \geq 0 \ t \geq 0 \ \varepsilon > 0$ **for** $s \ t \ \varepsilon$
<proof>

lemma *conv-in-prob*:
assumes $t \geq 0$
shows *tendsto-measure (proc-source X) X (X t) (at t within {0..})*
<proof>

lemma *conv-in-prob-finite*:
assumes $t \geq 0$
shows *tendsto-measure (proc-source X) X (X t) (at t within {0..T})*
<proof>

lemma *incr*: $\mathcal{P}(x \text{ in } \textit{source}. 2 \ \textit{powr} (-\gamma * n) \leq \textit{dist} (X ((k - 1) * 2 \ \textit{powr} - n) \ x) (X (k * 2 \ \textit{powr} - n) \ x))$
 $\leq C * 2 \ \textit{powr} (-n * (1+b-a*\gamma))$
if $k \geq 1 \ n \geq 0$ **for** $k \ n$
<proof>

end

In order to construct the modification of X , it suffices to construct a modification of X on $\{0..T\}$ for all finite T , from which we construct the modification on $\{0..\}$ via a countable union.

locale *Kolmogorov-Chentsov-finite* = *Kolmogorov-Chentsov* +
fixes $T :: \textit{real}$
assumes *zero-le-T*: $0 < T$
begin

A_n will characterise the set of states with increments that exceed the bounds required for Hölder continuity. As $n \rightarrow \infty$, this approaches the set of states for which X is not Hölder continuous. We define N as this limit, and show that N is a null set. On $\omega \in \Omega - N$, we show that $X(\omega)$ is Hölder continuous (and therefore uniformly continuous) on the dyadic rationals, and construct a modification by taking the continuous extension on the reals.

definition $A \equiv \lambda n. \textit{if} \ 2^{\wedge} n * T < 1 \ \textit{then} \ \textit{space} \ \textit{source} \ \textit{else}$

$\{x \in \text{space source.}$
 $\text{Max } \{ \text{dist } (X \text{ (real-of-int } (k - 1) * 2 \text{ powr} - \text{real } n) x) (X \text{ (real-of-int } k * 2$
 $\text{powr} - \text{real } n) x)$
 $\mid k. k \in \{1..[2^{\wedge}n * T]\} \} \geq 2 \text{ powr } (-\gamma * \text{real } n) \}$

abbreviation $B \equiv \lambda n. (\bigcup m. A (m + n))$

abbreviation $N \equiv \bigcap (\text{range } B)$

lemma *A-geq*: $2^{\wedge}n * T \geq 1 \implies A n = \{x \in \text{space source.}$

$\text{Max } \{ \text{dist } (X \text{ (real-of-int } (k - 1) * 2 \text{ powr} - \text{real } n) x) (X \text{ (real-of-int } k * 2$
 $\text{powr} - \text{real } n) x)$
 $\mid k. k \in \{1..[2^{\wedge}n * T]\} \} \geq 2 \text{ powr } (-\gamma * \text{real } n) \}$ **for** n
 $\langle \text{proof} \rangle$

lemma *A-measurable[measurable]*: $A n \in \text{sets source}$

$\langle \text{proof} \rangle$

lemma *emeasure-A-leg*:

fixes $n :: \text{nat}$

assumes [*simp*]: $2^{\wedge}n * T \geq 1$

shows $\text{emeasure source } (A n) \leq C * T * 2 \text{ powr } (-n * (b - a * \gamma))$

$\langle \text{proof} \rangle$

lemma *measure-A-leg*:

assumes $2^{\wedge}n * T \geq 1$

shows $\text{measure source } (A n) \leq C * T * 2 \text{ powr } (-n * (b - a * \gamma))$

$\langle \text{proof} \rangle$

lemma *summable-A*: $\text{summable } (\lambda m. \text{measure source } (A m))$

$\langle \text{proof} \rangle$

lemma *lim-B*: $(\lambda n. \text{measure source } (B n)) \longrightarrow 0$

$\langle \text{proof} \rangle$

lemma *N-null*: $N \in \text{null-sets source}$

$\langle \text{proof} \rangle$

lemma *notin-N-index*:

assumes $\omega \in \text{space source} - N$

obtains n_0 **where** $\omega \notin (\bigcup n. A (n + n_0))$

$\langle \text{proof} \rangle$

context

fixes ω

assumes $\omega: \omega \in \text{space source} - N$

begin

definition $n_0 \equiv \text{SOME } m. \omega \notin (\bigcup n. A (n + m)) \wedge m > 0$

lemma

shows $n\text{-zero}$: $\omega \notin (\bigcup n. A (n + n_0))$
and $n\text{-zero-nonzero}$: $n_0 > 0$

<proof>

lemma $n\text{zero-ge}$: $\bigwedge n. n \geq n_0 \implies 2^{\wedge n} * T \geq 1$

<proof>

lemma $\omega\text{-notin}$: $\bigwedge n. n \geq n_0 \implies \omega \notin A n$

<proof>

Klenke 21.7

lemma $X\text{-dyadic-incr}$:

assumes $k \in \{1..[2^{\wedge n} * T]\}$ $n \geq n_0$

shows $\text{dist } (X ((\text{real-of-int } k-1)/2^{\wedge n}) \omega) (X (\text{real-of-int } k/2^{\wedge n}) \omega) < 2^{\text{powr } (-\gamma * n)}$

<proof>

Klenke (21.8)

lemma dist-dyadic-mn :

assumes mn : $n_0 \leq n \leq m$

and $t\text{-dyadic}$: $t \in \text{dyadic-interval-step } m \ 0 \ T$

and $u\text{-dyadic-n}$: $u \in \text{dyadic-interval-step } n \ 0 \ T$

and ut : $u \leq t \ t - u < 2/2^{\wedge n}$

shows $\text{dist } (X u \omega) (X t \omega) \leq 2^{\text{powr } (-\gamma * n)} / (1 - 2^{\text{powr } -\gamma})$

<proof>

lemma dist-dyadic-fixed :

assumes mn : $n_0 \leq n \leq m$

and $s\text{-dyadic}$: $s \in \text{dyadic-interval-step } m \ 0 \ T$

and $t\text{-dyadic}$: $t \in \text{dyadic-interval-step } m \ 0 \ T$

and st : $s \leq t \ t - s \leq 1/2^{\wedge n}$

shows $\text{dist } (X t \omega) (X s \omega) \leq 2 * 2^{\text{powr } (-\gamma * n)} / (1 - 2^{\text{powr } -\gamma})$

<proof>

definition $C_0 \equiv 2 * 2^{\text{powr } \gamma} / (1 - 2^{\text{powr } -\gamma})$

lemma $C\text{-zero-ge}[simp]$: $C_0 > 0$

<proof>

Klenke (21.9)

Let $s, t \in D$ with $|s - t| \leq \frac{1}{2^n}$. By choosing the minimal $n \geq n_0$ such that $|t - s| \geq 2^{-n}$, we obtain by $\llbracket n_0 \leq ?n; ?n \leq ?m; ?s \in \text{dyadic-interval-step } ?m \ 0 \ T; ?t \in \text{dyadic-interval-step } ?m \ 0 \ T; ?s \leq ?t; ?t - ?s \leq 1 / 2^{?n} \rrbracket \implies \text{dist } (\text{process } X \ ?t \ \omega) (\text{process } X \ ?s \ \omega) \leq 2 * 2^{\text{powr } (-\gamma * \text{real } ?n)} / (1 - 2^{\text{powr } -\gamma})$:

$$|X_t(\omega) - X_s(\omega)| \leq C_0 |t - s|^\gamma$$

lemma *dist-dyadic*:

assumes $t: t \in \text{dyadic-interval } 0 \ T$

and $s: s \in \text{dyadic-interval } 0 \ T$

and *st-dist*: $\text{dist } t \ s \leq 1 / 2^{\wedge n_0}$

shows $\text{dist } (X \ t \ \omega) \ (X \ s \ \omega) \leq C_0 * (\text{dist } t \ s)^{\text{powr } \gamma}$

<proof>

definition $K \equiv C_0 * (2^{\wedge \text{nat } \lceil 2^{\wedge n_0} * T \rceil})^{\text{powr } (1 - \gamma)}$

lemma *C₀-le-K*: $C_0 \leq K$

<proof>

lemma *K-pos*: $0 < K$

<proof>

Klenke (21.10)

lemma *X-dyadic-le-K'*:

assumes *dyadic*: $s \in \text{dyadic-interval } 0 \ T \ t \in \text{dyadic-interval } 0 \ T$

and *st*: $s \leq t$

shows $\text{dist } (X \ s \ \omega) \ (X \ t \ \omega) \leq K * \text{dist } s \ t^{\text{powr } \gamma}$

<proof>

lemma *X-dyadic-le-K*:

assumes $s \in \text{dyadic-interval } 0 \ T$

and $t \in \text{dyadic-interval } 0 \ T$

shows $\text{dist } (X \ s \ \omega) \ (X \ t \ \omega) \leq K * \text{dist } s \ t^{\text{powr } \gamma}$

<proof>

corollary *holder-dyadic*: γ -holder-on (dyadic-interval 0 T) ($\lambda t. X \ t \ \omega$)

<proof>

lemma *uniformly-continuous-dyadic*: uniformly-continuous-on (dyadic-interval 0 T) ($\lambda t. X \ t \ \omega$)

<proof>

lemma *Lim-exists*: $\exists L. ((\lambda s. X \ s \ \omega) \longrightarrow L)$ (at t within (dyadic-interval 0 T))

if $t \in \{0..T\}$

<proof>

lemma *Lim-unique*: $\exists! L. ((\lambda s. X \ s \ \omega) \longrightarrow L)$ (at t within (dyadic-interval 0 T))

if $t \in \{0..T\}$

<proof>

definition $L \equiv (\lambda t. (\text{Lim } (\text{at } t \text{ within dyadic-interval } 0 \ T) (\lambda s. X \ s \ \omega)))$

lemma *X-tendsto-L*:

assumes $t \in \{0..T\}$

shows $((\lambda s. X s \omega) \longrightarrow L t)$ (at t within (dyadic-interval $0 T$))
 ⟨proof⟩

lemma *L-dist-K*:

assumes $s \in \{0..T\}$

and $t \in \{0..T\}$

shows $dist (L s) (L t) \leq K * dist s t \text{ powr } \gamma$

⟨proof⟩

corollary *L-holder: γ -holder-on $\{0..T\}$ L*

⟨proof⟩

corollary *L-local-holder: local-holder-on $\gamma \{0..T\}$ L*

⟨proof⟩

lemma *X-dyadic-eq-L*:

assumes $t \in \text{dyadic-interval } 0 T$

shows $X t \omega = L t$

⟨proof⟩

end

definition *default :: 'b where default = (SOME x. True)*

definition *X-tilde :: real \Rightarrow 'a \Rightarrow 'b where*

X-tilde \equiv ($\lambda t \omega$. if $\omega \in N$ then default else (Lim (at t within dyadic-interval $0 T$) ($\lambda s. X s \omega$)))

lemma *X-tilde-not-N-Lim*:

assumes $\omega \in \text{space source} - N$

shows $X\text{-tilde } t \omega = \text{Lim (at } t \text{ within dyadic-interval } 0 T) (\lambda s. X s \omega)$

⟨proof⟩

lemma *X-tilde-not-N-L*:

assumes $\omega \in \text{space source} - N$

shows $X\text{-tilde } t \omega = L \omega t$

⟨proof⟩

lemma *local-holder-X-tilde: local-holder-on $\gamma \{0..T\}$ ($\lambda t. X\text{-tilde } t \omega$)*

if $\omega \in \text{space source}$ **for** ω

⟨proof⟩

corollary *X-tilde-eq-L-AE: AE ω in source. X-tilde $t \omega = L \omega t$*

⟨proof⟩

corollary *X-tilde-eq-Lim-AE:*

AE ω in source. X-tilde $t \omega = \text{Lim (at } t \text{ within dyadic-interval } 0 T) (\lambda s. X s \omega)$

⟨proof⟩

lemma *X-tilde-tendsto-AE: $t \in \{0..T\} \implies \text{tendsto-AE source } X (X\text{-tilde } t)$ (at t*

within dyadic-interval $0 T$)
 ⟨proof⟩

end

context Kolmogorov-Chentsov-finite
begin

By (21.5) $0 \leq ?t \implies \text{tendsto-measure source (process } X \text{) (process } X \text{ } ?t \text{) (at } ?t \text{ within } \{0..?T\})$ and (21.11) $? \omega \in \text{space source} - (\bigcap_n \bigcup_m A (m + n)) \implies L ? \omega \equiv \lambda t. \text{Lim (at } t \text{ within dyadic-interval } 0 T \text{) } (\lambda s. \text{process } X \text{ } s ? \omega), P[X \neq \tilde{X}] = 0$

lemma *X-tilde-measurable[measurable]:*
assumes $t \in \{0..T\}$
shows *X-tilde* $t \in \text{borel-measurable source}$
 ⟨proof⟩

lemma *X-eq-X-tilde-AE: AE* ω *in source. X* $t \omega = \text{X-tilde } t \omega$ **if** $t \in \{0..T\}$ **for** t
 ⟨proof⟩

lemma *X-tilde-modification: modification (restrict-index X {0..T})*
(prob-space.process-of source (proc-target X) {0..T} X-tilde default)
 ⟨proof⟩

end

We have now shown that we can construct a modification of X for any interval $\{0..T\}$. We want to extend this result to construct a modification on the interval $\{0..\}$ - this can be constructed by gluing together all modifications with natural-valued T which results in a countable union of modifications, which itself is a modification.

context Kolmogorov-Chentsov
begin

lemma *Kolmogorov-Chentsov-finite: T > 0* \implies *Kolmogorov-Chentsov-finite X a b C* γT
 ⟨proof⟩

definition *Mod* $\equiv \lambda T. \text{SOME } Y. \text{modification (restrict-index X } \{0..T\}) Y \wedge (\forall x \in \text{space source. local-holder-on } \gamma \{0..T\} (\lambda t. Y t x))$

lemma *Mod: modification (restrict-index X {0..T}) (Mod T)*
($\forall x \in \text{space source. local-holder-on } \gamma \{0..T\} (\lambda t. (\text{Mod } T) t x)$) **if** $0 < T$ **for** T
 ⟨proof⟩

lemma *compatible-Mod: compatible (restrict-index X {0..T}) (Mod T)* **if** $0 < T$ **for** T
 ⟨proof⟩

lemma *Mod-source[simp]*: $\text{proc-source } (\text{Mod } T) = \text{source}$ **if** $0 < T$ **for** T
 ⟨proof⟩

lemma *Mod-target*: $\text{sets } (\text{proc-target } (\text{Mod } T)) = \text{sets } (\text{proc-target } X)$ **if** $0 < T$
for T
 ⟨proof⟩

lemma *Mod-index[simp]*: $0 < T \implies \text{proc-index } (\text{Mod } T) = \{0..T\}$
 ⟨proof⟩

lemma *indistinguishable-mod*:
indistinguishable ($\text{restrict-index } (\text{Mod } S) \{0..\text{min } S \ T\}$) ($\text{restrict-index } (\text{Mod } T)$
 $\{0..\text{min } S \ T\}$)
if $S > 0 \ T > 0$ **for** $S \ T$
 ⟨proof⟩

definition $N \ S \ T \equiv \text{SOME } N. N \in \text{null-sets source} \wedge \{\omega \in \text{space source}. \exists t \in$
 $\{0..\text{min } S \ T\}. (\text{Mod } S) \ t \ \omega \neq (\text{Mod } T) \ t \ \omega\} \subseteq N$

lemma N :
assumes $S > 0 \ T > 0$
shows $N \ S \ T \in \text{null-sets source} \wedge \{\omega \in \text{space source}. \exists t \in \{0..\text{min } S \ T\}. (\text{Mod}$
 $S) \ t \ \omega \neq (\text{Mod } T) \ t \ \omega\} \subseteq N \ S \ T$
 ⟨proof⟩

definition $N\text{-inf}$ **where** $N\text{-inf} \equiv (\bigcup S \in \mathbb{N} - \{0\}. (\bigcup T \in \mathbb{N} - \{0\}. N \ S \ T))$

lemma $N\text{-inf-null}$: $N\text{-inf} \in \text{null-sets source}$
 ⟨proof⟩

lemma *Mod-eq-N-inf*: $(\text{Mod } S) \ t \ \omega = (\text{Mod } T) \ t \ \omega$
if $\omega \in \text{space source} - N\text{-inf}$ $t \in \{0..\text{min } S \ T\}$ $S \in \mathbb{N} - \{0\}$ $T \in \mathbb{N} - \{0\}$ **for**
 $\omega \ t \ S \ T$
 ⟨proof⟩

definition $\text{default} :: 'b$ **where** $\text{default} = (\text{SOME } x. \text{True})$

definition $X\text{-mod} \equiv \lambda t \ \omega. \text{if } \omega \in \text{space source} - N\text{-inf} \text{ then } (\text{Mod } \lfloor t+1 \rfloor) \ t \ \omega \text{ else}$
 default

definition $X\text{-mod-process} \equiv \text{prob-space.process-of source } (\text{proc-target } X) \ \{0..\} \ X\text{-mod}$
 default

lemma *Mod-measurable[measurable]*: $\forall t \in \{0..\}. X\text{-mod } t \in \text{source} \rightarrow_M \text{proc-target}$
 X
 ⟨proof⟩

lemma *modification-X-mod-process*: $\text{modification } X \ X\text{-mod-process}$
 ⟨proof⟩

lemma *local-holder-X-mod*: *local-holder-on* $\gamma \{0..\}$ $(\lambda t. X\text{-mod } t \ \omega)$ **for** ω
 ⟨*proof*⟩

lemma *local-holder-X-mod-process*: *local-holder-on* $\gamma \{0..\}$ $(\lambda t. X\text{-mod-process } t \ \omega)$
for ω
 ⟨*proof*⟩

theorem *continuous-modification*:
 $\exists X'. \text{modification } X \ X' \wedge (\forall \omega. \text{local-holder-on } \gamma \{0..\} (\lambda t. X' \ t \ \omega))$
 ⟨*proof*⟩
end

theorem *Kolmogorov-Chentsov*:
fixes $X :: (\text{real}, 'a, 'b :: \text{polish-space}) \text{stochastic-process}$
and $a \ b \ C \ \gamma :: \text{real}$
assumes *index[simp]*: *proc-index* $X = \{0..\}$
and *target-borel[simp]*: *proc-target* $X = \text{borel}$
and *gt-0*: $a > 0 \ b > 0 \ C > 0$
and *b-leq-a*: $b \leq a$
and *gamma*: $\gamma \in \{0 < .. < b/a\}$
and *expectation*: $\bigwedge s \ t. \llbracket s \geq 0; t \geq 0 \rrbracket \implies$
 $(\int^+ x. \text{dist } (X \ t \ x) \ (X \ s \ x) \ \text{powr } a \ \partial \text{proc-source } X) \leq C * \text{dist } t \ s \ \text{powr}$
 $(1+b)$
shows $\exists X'. \text{modification } X \ X' \wedge (\forall \omega. \text{local-holder-on } \gamma \{0..\} (\lambda t. X' \ t \ \omega))$
 ⟨*proof*⟩
end

References

- [1] A. Klenke. *Probability theory: a comprehensive course*. Springer Science & Business Media, 2020.