

Formal Verification of an Explicit Counterexample to the Jacobian Conjecture

Arthur Freitas Ramos David Barros Hulak
Ruy J. G. B. de Queiroz

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Abstract

The Jacobian conjecture asks whether a polynomial self-map of affine space over a field of characteristic zero has a polynomial inverse whenever its Jacobian determinant is a nonzero constant [4, 5]. This entry gives an independent Isabelle/HOL verification of the explicit three-dimensional map announced by Levent Alpöge on July 20, 2026 [1]. The announcement credits Akhil Mathew with prompting the question and the AI system Claude Fable with work leading to the map; stable Lean and independent verification repositories appeared the same day [2, 3].

Isabelle proves that the formal Jacobian entries are the corresponding complex analytic partial derivatives and that the determinant is identically -2 . It verifies that three distinct rational points have the same image. Scaling the first output coordinate yields determinant exactly 1 while preserving noninjectivity. The theory defines polynomial invertibility and proves that the map has no polynomial inverse. It also constructs identity-padded polynomial maps, proves their block-Jacobian determinant is unchanged, and thereby obtains counterexamples in every finite dimension at least three. The development makes no claim about the two-dimensional conjecture.

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```

theory Jacobian-Counterexample
  imports
    Complex-Main
    HOL-Analysis.Derivative
    HOL-Analysis.Determinants
    HOL-Decision-Procs.Commutative-Ring
    HOL-Library.Cardinality
begin

```

1 Formal multivariate polynomials

We use a small expression datatype for multivariate polynomials. This makes polynomiality syntactic and gives a direct definition of formal partial differentiation, independent of analytic differentiability.

```

datatype ('v, 'a) poly-expr =
  PConst 'a
| PVar 'v
| PAdd ('v, 'a) poly-expr ('v, 'a) poly-expr
| PMul ('v, 'a) poly-expr ('v, 'a) poly-expr
| PPow ('v, 'a) poly-expr nat

instantiation poly-expr :: (type, zero) zero
begin

definition 0 = PConst 0

instance <proof>

end

fun eval-poly-expr ::
  ('v  $\Rightarrow$  'a::semiring-1)  $\Rightarrow$  ('v, 'a) poly-expr  $\Rightarrow$  'a where
  eval-poly-expr  $\alpha$  (PConst c) = c
| eval-poly-expr  $\alpha$  (PVar v) =  $\alpha$  v
| eval-poly-expr  $\alpha$  (PAdd p q) =
  eval-poly-expr  $\alpha$  p + eval-poly-expr  $\alpha$  q
| eval-poly-expr  $\alpha$  (PMul p q) =
  eval-poly-expr  $\alpha$  p * eval-poly-expr  $\alpha$  q
| eval-poly-expr  $\alpha$  (PPow p n) = eval-poly-expr  $\alpha$  p ^ n

fun partial-poly-expr ::
  'v  $\Rightarrow$  ('v, 'a::semiring-1) poly-expr  $\Rightarrow$  ('v, 'a) poly-expr where
  partial-poly-expr v (PConst c) = PConst 0
| partial-poly-expr v (PVar w) = PConst (if v = w then 1 else 0)
| partial-poly-expr v (PAdd p q) =
  PAdd (partial-poly-expr v p) (partial-poly-expr v q)
| partial-poly-expr v (PMul p q) =
  PAdd (PMul (partial-poly-expr v p) q)

```

$(PMul\ p\ (partial-poly-expr\ v\ q))$
 $| partial-poly-expr\ v\ (PPow\ p\ n) =$
 $PMul\ (PMul\ (PConst\ (of-nat\ n))\ (PPow\ p\ (n - 1)))$
 $(partial-poly-expr\ v\ p)$

lemma *eval-partial-poly-expr-has-field-derivative:*

fixes $\alpha :: 'v \Rightarrow complex$
shows $((\lambda u. eval-poly-expr\ (\alpha(v := u))\ p)\ has-field-derivative$
 $eval-poly-expr\ (\alpha(v := u))\ (partial-poly-expr\ v\ p))\ (at\ u)$
 $\langle proof \rangle$

definition *eval-poly-map ::*

$(('n, 'a::semiring-1)\ poly-expr) \wedge^n n \Rightarrow 'a \wedge^n n \Rightarrow 'a \wedge^n n$ **where**
 $eval-poly-map\ P\ x = (\chi\ i. eval-poly-expr\ (\lambda j. x\ \$\ j)\ (P\ \$\ i))$

definition *jacobian-matrix ::*

$(('n, 'a::comm-ring-1)\ poly-expr) \wedge^n n \Rightarrow 'a \wedge^n n \Rightarrow 'a \wedge^n n \wedge^n n$ **where**
 $jacobian-matrix\ P\ x =$
 $(\chi\ i\ j. eval-poly-expr\ (\lambda k. x\ \$\ k)\ (partial-poly-expr\ j\ (P\ \$\ i)))$

lemma *jacobian-matrix-entry-has-field-derivative:*

fixes $P :: (('n, complex)\ poly-expr) \wedge^n n$
shows $((\lambda u. eval-poly-expr\ ((\lambda k. x\ \$\ k)(j := u))\ (P\ \$\ i))$
 $has-field-derivative\ jacobian-matrix\ P\ x\ \$\ i\ \$\ j)\ (at\ (x\ \$\ j))$
 $\langle proof \rangle$

definition *polynomially-invertible ::*

$(('n, 'a::semiring-1)\ poly-expr) \wedge^n n \Rightarrow bool$ **where**
 $polynomially-invertible\ P \longleftrightarrow$
 $(\exists\ Q.$
 $(\forall x. eval-poly-map\ Q\ (eval-poly-map\ P\ x) = x) \wedge$
 $(\forall x. eval-poly-map\ P\ (eval-poly-map\ Q\ x) = x))$

lemma *polynomially-invertible-imp-inj:*

assumes $polynomially-invertible\ P$
shows $inj\ (eval-poly-map\ P)$
 $\langle proof \rangle$

definition *keller-map ::*

$(('n, 'a::comm-ring-1)\ poly-expr) \wedge^n n \Rightarrow bool$ **where**
 $keller-map\ P \longleftrightarrow$
 $(\exists\ d. d \neq 0 \wedge (\forall x. det\ (jacobian-matrix\ P\ x) = d))$

definition *normalized-keller-map ::*

$(('n, 'a::comm-ring-1)\ poly-expr) \wedge^n n \Rightarrow bool$ **where**
 $normalized-keller-map\ P \longleftrightarrow$
 $(\forall x. det\ (jacobian-matrix\ P\ x) = 1)$

definition *strong-keller-counterexample ::*

$((\text{'n}, \text{'a}::\text{comm-ring-1}) \text{ poly-expr}) \wedge \text{'n} \Rightarrow \text{bool}$ **where**
 $\text{strong-keller-counterexample } P \longleftrightarrow$
 $\text{keller-map } P \wedge \neg \text{inj } (\text{eval-poly-map } P)$

definition $\text{keller-counterexample} ::$
 $((\text{'n}, \text{'a}::\text{comm-ring-1}) \text{ poly-expr}) \wedge \text{'n} \Rightarrow \text{bool}$ **where**
 $\text{keller-counterexample } P \longleftrightarrow$
 $\text{keller-map } P \wedge \neg \text{polynomially-invertible } P$

definition $\text{strong-normalized-keller-counterexample} ::$
 $((\text{'n}, \text{'a}::\text{comm-ring-1}) \text{ poly-expr}) \wedge \text{'n} \Rightarrow \text{bool}$ **where**
 $\text{strong-normalized-keller-counterexample } P \longleftrightarrow$
 $\text{normalized-keller-map } P \wedge \neg \text{inj } (\text{eval-poly-map } P)$

definition $\text{normalized-keller-counterexample} ::$
 $((\text{'n}, \text{'a}::\text{comm-ring-1}) \text{ poly-expr}) \wedge \text{'n} \Rightarrow \text{bool}$ **where**
 $\text{normalized-keller-counterexample } P \longleftrightarrow$
 $\text{normalized-keller-map } P \wedge \neg \text{polynomially-invertible } P$

lemma $\text{strong-keller-counterexample-imp-keller-counterexample}:$
 $\text{strong-keller-counterexample } P \implies \text{keller-counterexample } P$
 $\langle \text{proof} \rangle$

lemma $\text{strong-normalized-keller-counterexample-imp-counterexample}:$
 $\text{strong-normalized-keller-counterexample } P \implies$
 $\text{normalized-keller-counterexample } P$
 $\langle \text{proof} \rangle$

2 The announced polynomial map

Alpöge announced the following three-coordinate polynomial map on July 20, 2026 [1]. The announcement credits Akhil Mathew with prompting the question and the AI system Claude Fable with work leading to the example; stable independent verification repositories appeared the same day [2, 3]. The definitions below encode the map syntactically, so the subsequent Jacobian calculation is a theorem about formal polynomial differentiation rather than an imported computer-algebra result.

definition $x\text{-poly} :: (\mathcal{A}, \text{'a}::\text{comm-ring-1}) \text{ poly-expr}$ **where**
 $x\text{-poly} = PVar\ 1$

definition $y\text{-poly} :: (\mathcal{A}, \text{'a}::\text{comm-ring-1}) \text{ poly-expr}$ **where**
 $y\text{-poly} = PVar\ 2$

definition $z\text{-poly} :: (\mathcal{A}, \text{'a}::\text{comm-ring-1}) \text{ poly-expr}$ **where**
 $z\text{-poly} = PVar\ 3$

definition $\text{one-plus-xy-poly} :: (\mathcal{A}, \text{'a}::\text{comm-ring-1}) \text{ poly-expr}$ **where**

$one-plus-xy-poly = PAdd (PConst 1) (PMul x-poly y-poly)$

definition $four-plus-three-xy-poly :: (\mathcal{A}, 'a::comm-ring-1) poly-expr$ **where**
 $four-plus-three-xy-poly =$
 $PAdd (PConst 4) (PMul (PConst 3) (PMul x-poly y-poly))$

definition $component-a-poly :: (\mathcal{A}, 'a::comm-ring-1) poly-expr$ **where**
 $component-a-poly =$
 $PAdd (PMul (PPow one-plus-xy-poly 3) z-poly)$
 $(PMul (PMul (PPow y-poly 2) one-plus-xy-poly)$
 $four-plus-three-xy-poly))$

definition $component-b-poly :: (\mathcal{A}, 'a::comm-ring-1) poly-expr$ **where**
 $component-b-poly =$
 $PAdd y-poly$
 $(PAdd (PMul (PMul (PConst 3) x-poly)$
 $(PMul (PPow one-plus-xy-poly 2) z-poly))$
 $(PMul (PMul (PMul (PConst 3) x-poly) (PPow y-poly 2))$
 $four-plus-three-xy-poly)))$

definition $component-c-poly :: (\mathcal{A}, 'a::comm-ring-1) poly-expr$ **where**
 $component-c-poly =$
 $PAdd (PMul (PConst 2) x-poly)$
 $(PAdd (PMul (PConst (-3)) (PMul (PPow x-poly 2) y-poly))$
 $(PMul (PConst (-1)) (PMul (PPow x-poly 3) z-poly)))$

definition $alpoge-map :: ((\mathcal{A}, 'a::comm-ring-1) poly-expr)^{\mathcal{A}}$ **where**
 $alpoge-map =$
 $vector [component-a-poly, component-b-poly, component-c-poly]$

lemma $eval-alpoge-map$:
 $eval-poly-map alpoge-map (vector [x, y, z]) =$
 $vector [$
 $(1 + x * y)^{\mathcal{A}} * z + y^2 * (1 + x * y) * (4 + 3 * x * y),$
 $y + 3 * x * (1 + x * y)^2 * z + 3 * x * y^2 * (4 + 3 * x * y),$
 $2 * x - 3 * x^2 * y - x^3 * z]$
 $\langle proof \rangle$

corollary $alpoge-jacobian-entry-has-field-derivative$:
fixes $x :: complex^{\mathcal{A}}$
shows $((\lambda u.$
 $eval-poly-map (alpoge-map :: ((\mathcal{A}, complex) poly-expr)^{\mathcal{A}})$
 $(\chi k. if k = j then u else x \$ k) \$ i)$
 $has-field-derivative$
 $jacobian-matrix (alpoge-map :: ((\mathcal{A}, complex) poly-expr)^{\mathcal{A}}) x \$ i \$ j)$
 $(at (x \$ j))$
 $\langle proof \rangle$

lemma $jacobian-alpoge-map$:

jacobian-matrix alpoge-map (*vector* [*x*, *y*, *z*]) =
vector [
vector [
 $3 * y * (1 + x * y)^2 * z +$
 $y^3 * (4 + 3 * x * y) + 3 * y^3 * (1 + x * y),$
 $3 * x * (1 + x * y)^2 * z +$
 $2 * y * (1 + x * y) * (4 + 3 * x * y) +$
 $x * y^2 * (4 + 3 * x * y) + 3 * x * y^2 * (1 + x * y),$
 $(1 + x * y)^3],$
vector [
 $3 * (1 + x * y)^2 * z + 6 * x * y * (1 + x * y) * z +$
 $3 * y^2 * (4 + 3 * x * y) + 9 * x * y^3,$
 $1 + 6 * x^2 * (1 + x * y) * z +$
 $6 * x * y * (4 + 3 * x * y) + 9 * x^2 * y^2,$
 $3 * x * (1 + x * y)^2],$
vector [
 $2 - 6 * x * y - 3 * x^2 * z,$
 $-3 * x^2,$
 $-(x^3)]]$
 ⟨*proof*⟩

lemma *alpoge-jacobian-det-coordinates:*

det (*jacobian-matrix alpoge-map* (*vector* [*x*, *y*, *z*])) = -2
 ⟨*proof*⟩

theorem *alpoge-jacobian-det:*

det (*jacobian-matrix alpoge-map* *p*) = -2
 ⟨*proof*⟩

3 An explicit three-point fibre

definition *collision-point-0* :: 'a::field-char-0^3 **where**

collision-point-0 = *vector* [0, 0, -1 / 4]

definition *collision-point-pos* :: 'a::field-char-0^3 **where**

collision-point-pos = *vector* [1, -3 / 2, 13 / 2]

definition *collision-point-neg* :: 'a::field-char-0^3 **where**

collision-point-neg = *vector* [-1, 3 / 2, 13 / 2]

definition *collision-value* :: 'a::field-char-0^3 **where**

collision-value = *vector* [-1 / 4, 0, 0]

lemma *collision-points-distinct:*

distinct [*collision-point-0*, *collision-point-pos*, *collision-point-neg*]
 ⟨*proof*⟩

lemma *alpoge-collision:*

eval-poly-map alpoge-map collision-point-0 = *collision-value*

$eval_poly_map \ alpoge_map \ collision_point_pos = collision_value$
 $eval_poly_map \ alpoge_map \ collision_point_neg = collision_value$
 $\langle proof \rangle$

lemma *alpoge-map-not-injective:*

$\neg inj$
 $(eval_poly_map \ (alpoge_map :: ((\mathbb{Z}, 'a::field_char\ 0) \ poly_expr)^\mathbb{Z}))$
 $\langle proof \rangle$

theorem *alpoge-map-is-strong-keller-counterexample:*

$strong_keller_counterexample$
 $(alpoge_map :: ((\mathbb{Z}, 'a::field_char\ 0) \ poly_expr)^\mathbb{Z})$
 $\langle proof \rangle$

corollary *alpoge-map-is-keller-counterexample:*

$keller_counterexample$
 $(alpoge_map :: ((\mathbb{Z}, 'a::field_char\ 0) \ poly_expr)^\mathbb{Z})$
 $\langle proof \rangle$

corollary *alpoge-map-not-polynomially-invertible:*

$\neg polynomially_invertible$
 $(alpoge_map :: ((\mathbb{Z}, 'a::field_char\ 0) \ poly_expr)^\mathbb{Z})$
 $\langle proof \rangle$

4 A determinant-one normalization

definition *normalized-alpoge-map ::*

$((\mathbb{Z}, 'a::field_char\ 0) \ poly_expr)^\mathbb{Z}$ **where**
 $normalized_alpoge_map =$
 $vector \ [$
 $\quad PMul \ (PConst \ (-1 \ / \ 2)) \ component_a_poly,$
 $\quad component_b_poly,$
 $\quad component_c_poly]$

lemma *eval-normalized-alpoge-map:*

$eval_poly_map \ normalized_alpoge_map \ (vector \ [x, y, z]) =$
 $vector \ [$
 $\quad (-1 \ / \ 2) * ((1 + x * y)^\mathbb{Z} * z + y^\mathbb{Z} * (1 + x * y) * (4 + \mathbb{Z} * x * y)),$
 $\quad y + \mathbb{Z} * x * (1 + x * y)^\mathbb{Z} * z + \mathbb{Z} * x * y^\mathbb{Z} * (4 + \mathbb{Z} * x * y),$
 $\quad 2 * x - \mathbb{Z} * x^\mathbb{Z} * y - x^\mathbb{Z} * z]$
 $\langle proof \rangle$

lemma *jacobian-normalized-alpoge-map:*

$jacobian_matrix \ normalized_alpoge_map \ p =$
 $(\chi \ i \ j.$
 $\quad if \ i = 1 \ then$
 $\quad \quad (-1 \ / \ 2) * (jacobian_matrix \ alpoge_map \ p \ \$ \ i \ \$ \ j)$
 $\quad else \ jacobian_matrix \ alpoge_map \ p \ \$ \ i \ \$ \ j)$

<proof>

theorem *normalized- alpoge -jacobian-det:*

$\det (\text{jacobian-matrix } \text{normalized-}\text{alpoge-map } p) = 1$

<proof>

lemma *normalized- alpoge -collision:*

$\text{eval-poly-map } \text{normalized-}\text{alpoge-map } \text{collision-point-0} =$

$\text{vector } [1 / 8, 0, 0]$

$\text{eval-poly-map } \text{normalized-}\text{alpoge-map } \text{collision-point-pos} =$

$\text{vector } [1 / 8, 0, 0]$

$\text{eval-poly-map } \text{normalized-}\text{alpoge-map } \text{collision-point-neg} =$

$\text{vector } [1 / 8, 0, 0]$

<proof>

lemma *normalized- alpoge -map-not-injective:*

$\neg \text{inj}$

$(\text{eval-poly-map}$

$(\text{normalized-}\text{alpoge-map} :: ((\mathbb{F}, 'a::\text{field-char-0}) \text{poly-expr})^{\wedge 3}))$

<proof>

theorem *normalized- alpoge -map-is-strong-counterexample:*

strong-normalized-keller-counterexample

$(\text{normalized-}\text{alpoge-map} :: ((\mathbb{F}, 'a::\text{field-char-0}) \text{poly-expr})^{\wedge 3})$

<proof>

theorem *normalized- alpoge -map-is-counterexample:*

normalized-keller-counterexample

$(\text{normalized-}\text{alpoge-map} :: ((\mathbb{F}, 'a::\text{field-char-0}) \text{poly-expr})^{\wedge 3})$

<proof>

corollary *normalized- alpoge -map-not-polynomially-invertible:*

$\neg \text{polynomially-invertible}$

$(\text{normalized-}\text{alpoge-map} :: ((\mathbb{F}, 'a::\text{field-char-0}) \text{poly-expr})^{\wedge 3})$

<proof>

This is the determinant-one form of the counterexample. Since every coordinate is represented by a polynomial expression, the theorem directly refutes the normalized dimension-three Jacobian conjecture over any field of characteristic zero.

corollary *complex-jacobian-conjecture-counterexample:*

normalized-keller-counterexample

$(\text{normalized-}\text{alpoge-map} :: ((\mathbb{F}, \text{complex}) \text{poly-expr})^{\wedge 3})$

<proof>

5 Identity padding in higher dimensions

lemma *prod-UNIV-sum:*

fixes $f :: ('n::finite + 'm::finite) \Rightarrow 'a::comm-monoid-mult$
shows $(\prod_{x \in UNIV}. f\ x) =$
 $(\prod_{i \in UNIV}. f\ (Inl\ i)) * (\prod_{j \in UNIV}. f\ (Inr\ j))$
 $\langle proof \rangle$

definition *block-identity-matrix* ::
 $'a::semiring-1 \wedge 'n \wedge 'n \Rightarrow 'a \wedge ('n + 'm::finite) \wedge ('n + 'm)$ **where**
block-identity-matrix $A =$
 $(\chi\ i\ j.$
 $\text{case } i \text{ of}$
 $Inl\ r \Rightarrow (\text{case } j \text{ of } Inl\ c \Rightarrow A\ \$\ r\ \$\ c \mid Inr\ - \Rightarrow 0)$
 $\mid Inr\ r \Rightarrow (\text{case } j \text{ of } Inl\ - \Rightarrow 0 \mid Inr\ c \Rightarrow \text{if } r = c \text{ then } 1 \text{ else } 0))$

lemma *det-block-identity-matrix*:
fixes $A :: 'a::comm-ring-1 \wedge 'n::finite \wedge 'n$
shows $\det\ (\text{block-identity-matrix } A :: 'a \wedge ('n + 'm::finite) \wedge ('n + 'm)) = \det\ A$
 $\langle proof \rangle$

fun *rename-poly-expr* ::
 $('v \Rightarrow 'w) \Rightarrow ('v, 'a)\ \text{poly-expr} \Rightarrow ('w, 'a)\ \text{poly-expr}$ **where**
 $\text{rename-poly-expr } f\ (PConst\ c) = PConst\ c$
 $\mid \text{rename-poly-expr } f\ (PVar\ v) = PVar\ (f\ v)$
 $\mid \text{rename-poly-expr } f\ (PAdd\ p\ q) =$
 $PAdd\ (\text{rename-poly-expr } f\ p)\ (\text{rename-poly-expr } f\ q)$
 $\mid \text{rename-poly-expr } f\ (PMul\ p\ q) =$
 $PMul\ (\text{rename-poly-expr } f\ p)\ (\text{rename-poly-expr } f\ q)$
 $\mid \text{rename-poly-expr } f\ (PPow\ p\ n) = PPow\ (\text{rename-poly-expr } f\ p)\ n$

lemma *eval-rename-poly-expr*:
 $\text{eval-poly-expr } \alpha\ (\text{rename-poly-expr } f\ p) =$
 $\text{eval-poly-expr } (\alpha \circ f)\ p$
 $\langle proof \rangle$

lemma *partial-rename-poly-expr-Inl*:
 $\text{partial-poly-expr } (Inl\ v)\ (\text{rename-poly-expr } Inl\ p) =$
 $\text{rename-poly-expr } Inl\ (\text{partial-poly-expr } v\ p)$
 $\langle proof \rangle$

lemma *eval-partial-rename-poly-expr-Inr*:
 $\text{eval-poly-expr } \alpha$
 $(\text{partial-poly-expr } (Inr\ w)\ (\text{rename-poly-expr } Inl\ p)) = 0$
 $\langle proof \rangle$

definition *join-vector* ::
 $'a \wedge 'n \Rightarrow 'a \wedge 'm \Rightarrow 'a \wedge ('n + 'm)$ **where**
 $\text{join-vector } x\ y =$
 $(\chi\ i. \text{case } i \text{ of } Inl\ j \Rightarrow x\ \$\ j \mid Inr\ j \Rightarrow y\ \$\ j)$

lemma *join-vector-Inl [simp]*: $\text{join-vector } x\ y\ \$\ Inl\ i = x\ \$\ i$

$\langle \text{proof} \rangle$

lemma *join-vector-Inr* [simp]: $\text{join-vector } x \ y \ \$ \text{ Inr } j = y \ \$ \ j$
 $\langle \text{proof} \rangle$

lemma *join-vector-comp-Inl* [simp]:
 $(\$) (\text{join-vector } x \ y) \circ \text{Inl} = (\$) \ x$
 $\langle \text{proof} \rangle$

lemma *join-vector-comp-Inr* [simp]:
 $(\$) (\text{join-vector } x \ y) \circ \text{Inr} = (\$) \ y$
 $\langle \text{proof} \rangle$

lemma *join-vector-eq-iff* [simp]:
 $\text{join-vector } x \ y = \text{join-vector } x' \ y' \iff x = x' \wedge y = y'$
 $\langle \text{proof} \rangle$

definition *pad-poly-map* ::
 $((n, 'a::\text{semiring-1}) \text{ poly-expr})^\wedge n \Rightarrow$
 $((n + 'm::\text{finite}, 'a) \text{ poly-expr})^\wedge (n + 'm) \text{ where}$
 $\text{pad-poly-map } P =$
 $(\chi \ i.$
 $\text{case } i \text{ of}$
 $\text{Inl } j \Rightarrow \text{rename-poly-expr Inl } (P \ \$ \ j)$
 $| \text{Inr } j \Rightarrow \text{PVar } (\text{Inr } j))$

lemma *eval-pad-poly-map*:
 $\text{eval-poly-map } (\text{pad-poly-map } P) (\text{join-vector } x \ y) =$
 $\text{join-vector } (\text{eval-poly-map } P \ x) \ y$
 $\langle \text{proof} \rangle$

lemma *jacobian-pad-poly-map*:
 $\text{jacobian-matrix } (\text{pad-poly-map } P) (\text{join-vector } x \ y) =$
 $\text{block-identity-matrix } (\text{jacobian-matrix } P \ x)$
 $\langle \text{proof} \rangle$

corollary *jacobian-det-pad-poly-map*:
 $\det (\text{jacobian-matrix } (\text{pad-poly-map } P) (\text{join-vector } x \ y)) =$
 $\det (\text{jacobian-matrix } P \ x)$
 $\langle \text{proof} \rangle$

definition *left-subvector* :: $'a^\wedge (n::\text{finite} + 'm::\text{finite}) \Rightarrow 'a^\wedge n$ **where**
 $\text{left-subvector } z = (\chi \ i. z \ \$ \ \text{Inl } i)$

definition *right-subvector* :: $'a^\wedge (n::\text{finite} + 'm::\text{finite}) \Rightarrow 'a^\wedge m$ **where**
 $\text{right-subvector } z = (\chi \ j. z \ \$ \ \text{Inr } j)$

lemma *join-vector-subvectors* [simp]:
 $\text{join-vector } (\text{left-subvector } z) (\text{right-subvector } z) = z$

<proof>

lemma *pad-poly-map-not-injective:*
fixes $P :: (('n, 'a::semiring-1) \text{poly-expr}) \frown 'n$
assumes $\neg \text{inj } (\text{eval-poly-map } P)$
shows $\neg \text{inj}$
 $(\text{eval-poly-map}$
 $(\text{pad-poly-map } P ::$
 $((('n + 'm::finite, 'a) \text{poly-expr}) \frown ('n + 'm))))$
<proof>

lemma *strong-normalized-keller-counterexample-pad-poly-map:*
fixes $P :: (('n, 'a::comm-ring-1) \text{poly-expr}) \frown 'n$
assumes *strong-normalized-keller-counterexample* P
shows *strong-normalized-keller-counterexample*
 $(\text{pad-poly-map } P ::$
 $((('n + 'm::finite, 'a) \text{poly-expr}) \frown ('n + 'm)))$
<proof>

theorem *normalized-alpoge-padding-is-strong-counterexample:*
strong-normalized-keller-counterexample
 $(\text{pad-poly-map}$
 $(\text{normalized-alpoge-map} :: ((\mathfrak{Z}, 'a::field-char-0) \text{poly-expr}) \frown \mathfrak{Z}) ::$
 $((\mathfrak{Z} + 'm::finite, 'a) \text{poly-expr}) \frown (\mathfrak{Z} + 'm))$
<proof>

theorem *normalized-alpoge-padding-is-counterexample:*
normalized-keller-counterexample
 $(\text{pad-poly-map}$
 $(\text{normalized-alpoge-map} :: ((\mathfrak{Z}, 'a::field-char-0) \text{poly-expr}) \frown \mathfrak{Z}) ::$
 $((\mathfrak{Z} + 'm::finite, 'a) \text{poly-expr}) \frown (\mathfrak{Z} + 'm))$
<proof>

corollary *complex-jacobian-conjecture-counterexample-in-higher-dimension:*
normalized-keller-counterexample
 $(\text{pad-poly-map}$
 $(\text{normalized-alpoge-map} :: ((\mathfrak{Z}, \text{complex}) \text{poly-expr}) \frown \mathfrak{Z}) ::$
 $((\mathfrak{Z} + 'm::finite, \text{complex}) \text{poly-expr}) \frown (\mathfrak{Z} + 'm))$
<proof>

lemma *padded-dimension-card:*
 $\text{CARD}(\mathfrak{Z} + 'm::finite) = \mathfrak{Z} + \text{CARD}('m)$
<proof>

lemma *padded-dimension-greater-than-three:*
 $\mathfrak{Z} < \text{CARD}(\mathfrak{Z} + 'm::finite)$
<proof>

The extra finite type $'m$ is nonempty, so the padded coordinate type has dimension strictly greater than three. Conversely, every concrete finite

dimension above three can be represented by choosing an extra finite type of the required cardinality. Thus the dimension-three counterexample yields counterexamples in all dimensions $n > 3$. Together with the original theorem this covers every $n \geq 3$. Nothing in this development addresses the two-dimensional conjecture.

end

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