

From Abstract to Concrete Gödel’s Incompleteness Theorems—Part I

Andrei Popescu Dmitriy Traytel

February 6, 2026

Abstract

We validate an abstract formulation of Gödel’s First and Second Incompleteness Theorems from a [separate AFP entry](#) by instantiating them to the case of *finite sound extensions of the Hereditarily Finite (HF) Set theory*, i.e., FOL theories extending the HF Set theory with a finite set of axioms that are sound in the standard model. The concrete results had been previously formalised in an [AFP entry by Larry Paulson](#); our instantiation reuses the infrastructure developed in that entry.

Contents

1 The Instantiation

1

1 The Instantiation

definition $Fvars\ t = \{a :: name. \neg\ atom\ a\ \#\ t\}$

lemma $Fvars_tm_simps[simp]$:

$Fvars\ Zero = \{\}$
 $Fvars\ (Var\ a) = \{a\}$
 $Fvars\ (Eats\ x\ y) = Fvars\ x \cup Fvars\ y$
by (*auto simp: Fvars_def fresh_at_base(2)*)

lemma $finite_Fvars_tm[simp]$:

fixes $t :: tm$
shows $finite\ (Fvars\ t)$
by (*induct t rule: tm.induct*) *auto*

lemma $Fvars_fm_simps[simp]$:

$Fvars\ (x\ IN\ y) = Fvars\ x \cup Fvars\ y$
 $Fvars\ (x\ EQ\ y) = Fvars\ x \cup Fvars\ y$
 $Fvars\ (A\ OR\ B) = Fvars\ A \cup Fvars\ B$
 $Fvars\ (A\ AND\ B) = Fvars\ A \cup Fvars\ B$
 $Fvars\ (A\ IMP\ B) = Fvars\ A \cup Fvars\ B$
 $Fvars\ Fls = \{\}$
 $Fvars\ (Neg\ A) = Fvars\ A$
 $Fvars\ (Ex\ a\ A) = Fvars\ A - \{a\}$
 $Fvars\ (All\ a\ A) = Fvars\ A - \{a\}$
by (*auto simp: Fvars_def fresh_at_base(2)*)

lemma $finite_Fvars_fm[simp]$:

fixes $A :: fm$
shows $finite\ (Fvars\ A)$

by (induct A rule: fm.induct) auto

lemma subst_tm_subst_tm[simp]:

$x \neq y \implies \text{atom } x \# u \implies \text{subst } y \ u \ (\text{subst } x \ t \ v) = \text{subst } x \ (\text{subst } y \ u \ t) \ (\text{subst } y \ u \ v)$

by (induct v rule: tm.induct) auto

lemma subst_fm_subst_fm[simp]:

$x \neq y \implies \text{atom } x \# u \implies (A(x::=t))(y:=u) = (A(y:=u))(x::=\text{subst } y \ u \ t)$

by (nominal_induct A avoiding: x t y u rule: fm.strong_induct) auto

lemma Fvars_ground_aux: $Fvars \ t \subseteq B \implies \text{ground_aux } t \ (\text{atom } 'B)$

by (induct t rule: tm.induct) auto

lemma ground_Fvars: $\text{ground } t \longleftrightarrow Fvars \ t = \{\}$

apply (rule iffI)

apply (auto simp only: Fvars_def ground_fresh) []

apply (auto intro: Fvars_ground_aux[of t {}], simplified)

done

lemma Fvars_ground_fm_aux: $Fvars \ A \subseteq B \implies \text{ground_fm_aux } A \ (\text{atom } 'B)$

apply (induct A arbitrary: B rule: fm.induct)

apply (auto simp: Diff_subset_conv Fvars_ground_aux)

apply (drule meta_spec, drule meta_mp, assumption)

apply auto

done

lemma ground_fm_Fvars: $\text{ground_fm } A \longleftrightarrow Fvars \ A = \{\}$

apply (rule iffI)

apply (auto simp only: Fvars_def ground_fresh) []

apply (auto intro: Fvars_ground_fm_aux[of A {}], simplified)

done

interpretation Generic_Syntax where

var = UNIV :: name set

and trm = UNIV :: tm set

and fmla = UNIV :: fm set

and Var = Var

and FvarsT = Fvars

and substT = $\lambda t \ u \ x. \text{subst } x \ u \ t$

and Fvars = Fvars

and subst = $\lambda A \ u \ x. \text{subst_fm } A \ x \ u$

apply unfold_locales

subgoal by simp

subgoal by simp

subgoal by simp

subgoal by simp

subgoal by simp

subgoal by simp

subgoal by simp

subgoal by simp

subgoal for t by (induct t rule: tm.induct) auto

subgoal by simp

subgoal by simp

subgoal by simp

subgoal unfolding Fvars_def fresh_subst_fm_if by auto

subgoal unfolding Fvars_def by auto

subgoal unfolding Fvars_def by simp

subgoal by simp

subgoal unfolding *Fvars_def* **by** *simp*
done

lemma *coding_tm_Fvars_empty*[*simp*]: *coding_tm t* \implies *Fvars t* = {}
by (*induct t rule: coding_tm.induct*) (*auto simp: Fvars_def*)

lemma *Fvars_empty_ground*[*simp*]: *Fvars t* = {} \implies *ground t*
by (*induct t rule: tm.induct*) *auto*

interpretation *Syntax_with_Numerals* **where**

var = *UNIV* :: *name set*
and *trm* = *UNIV* :: *tm set*
and *fmla* = *UNIV* :: *fm set*
and *num* = {*t. ground t*}
and *Var* = *Var*
and *FvarsT* = *Fvars*
and *substT* = $\lambda t u x. \text{subst } x u t$
and *Fvars* = *Fvars*
and *subst* = $\lambda A u x. \text{subst_fm } A x u$
apply *unfold_locales*
subgoal by (*auto intro!: exI[of _ Zero]*)
subgoal by *simp*
subgoal by (*simp add: ground_Fvars*)
done

declare *FvarsT_num*[*simp del*]

interpretation *Deduct2_with_False* **where**

var = *UNIV* :: *name set*
and *trm* = *UNIV* :: *tm set*
and *fmla* = *UNIV* :: *fm set*
and *num* = {*t. ground t*}
and *Var* = *Var*
and *FvarsT* = *Fvars*
and *substT* = $\lambda t u x. \text{subst } x u t$
and *Fvars* = *Fvars*
and *subst* = $\lambda A u x. \text{subst_fm } A x u$
and *eql* = (*EQ*)
and *cnj* = (*AND*)
and *imp* = (*IMP*)
and *all* = *All*
and *exi* = *Ex*
and *fls* = *Fls*
and *prv* = (\vdash) {}
and *bprv* = (\vdash) {}
apply *unfold_locales*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*
subgoal by *simp*

```

subgoal unfolding Fvars_def by simp
subgoal unfolding Fvars_def by simp
subgoal by simp
subgoal by simp
subgoal by simp
subgoal by simp
subgoal using MP_null by blast
subgoal by blast
subgoal for A B C
  apply (rule Imp_I)+
  apply (rule MP_same[of _ B])
  apply (rule MP_same[of _ C])
  apply (auto intro: Neg_D)
done
subgoal by blast
subgoal by blast
subgoal by blast
subgoal unfolding Fvars_def by (auto intro: MP_null)
subgoal unfolding Fvars_def by (auto intro: MP_null)
subgoal by (auto intro: All_D)
subgoal by (auto intro: Ex_I)
subgoal by simp
subgoal by (metis Conj_E2 Iff_def Imp_I Var_Eq_subst_Iff)
subgoal by blast
subgoal by simp
done

```

interpretation HBL1 where

```

  var = UNIV :: name set
and trm = UNIV :: tm set
and fmla = UNIV :: fm set
and num = {t. ground t}
and Var = Var
and FvarsT = Fvars
and substT =  $\lambda t u x. \text{subst } x u t$ 
and Fvars = Fvars
and subst =  $\lambda A u x. \text{subst\_fm } A x u$ 
and eql = (EQ)
and cnj = (AND)
and imp = (IMP)
and all = All
and exi = Ex
and prv = ( $\vdash$ ) {}
and bprv = ( $\vdash$ ) {}
and enc = quot
and P = PfP (Var xx)
apply unfold_locales
subgoal by (simp add: quot_fm_coding)
subgoal by simp
subgoal unfolding Fvars_def by (auto simp: fresh_at_base(2))
subgoal by (auto simp: proved_imp_proved_PfP)
done

```

interpretation Goedel_Form where

```

  var = UNIV :: name set
and trm = UNIV :: tm set
and fmla = UNIV :: fm set
and num = {t. ground t}

```

```

and  $Var = Var$ 
and  $FvarsT = Fvars$ 
and  $substT = \lambda t u x. subst\ x\ u\ t$ 
and  $Fvars = Fvars$ 
and  $subst = \lambda A u x. subst\_fm\ A\ x\ u$ 
and  $eql = (EQ)$ 
and  $cnj = (AND)$ 
and  $imp = (IMP)$ 
and  $all = All$ 
and  $exi = Ex$ 
and  $fls = Fls$ 
and  $prv = (\vdash)\ \{\}$ 
and  $bprv = (\vdash)\ \{\}$ 
and  $enc = quot$ 
and  $S = KRP\ (quot\ (Var\ xx))\ (Var\ xx)\ (Var\ yy)$ 
and  $P = Pfp\ (Var\ xx)$ 
apply  $unfold\_locales$ 
subgoal by  $simp$ 
subgoal unfolding  $Fvars\_def$  by  $(auto\ simp: fresh\_at\_base(2))$ 
subgoal
  unfolding  $Let\_def$ 
  by  $(subst\ psubst\_eq\_rawpsubst2)$ 
   $(auto\ simp: quot\_fm\_coding\ prove\_KRP\ Fvars\_def)$ 
subgoal
  unfolding  $Let\_def$ 
  by  $(subst\ (1\ 2)\ psubst\_eq\_rawpsubst2)$ 
   $(auto\ simp: quot\_fm\_coding\ KRP\_unique[THEN\ Sym]\ Fvars\_def)$ 
done

```

interpretation $g2: Goedel_Second_Assumptions$ **where**

```

   $var = UNIV :: name\ set$ 
and  $trm = UNIV :: tm\ set$ 
and  $fmla = UNIV :: fm\ set$ 
and  $num = \{t. ground\ t\}$ 
and  $Var = Var$ 
and  $FvarsT = Fvars$ 
and  $substT = \lambda t u x. subst\ x\ u\ t$ 
and  $Fvars = Fvars$ 
and  $subst = \lambda A u x. subst\_fm\ A\ x\ u$ 
and  $eql = (EQ)$ 
and  $cnj = (AND)$ 
and  $imp = (IMP)$ 
and  $all = All$ 
and  $exi = Ex$ 
and  $fls = Fls$ 
and  $prv = (\vdash)\ \{\}$ 
and  $bprv = (\vdash)\ \{\}$ 
and  $enc = quot$ 
and  $S = KRP\ (quot\ (Var\ xx))\ (Var\ xx)\ (Var\ yy)$ 
and  $P = Pfp\ (Var\ xx)$ 
apply  $unfold\_locales$ 
subgoal by  $(auto\ simp: PP\_def\ intro: Pfp\_implies\_ModPon\_Pfp\_quot)$ 
subgoal by  $(auto\ simp: PP\_def\ quot\_fm\_coding\ Provability)$ 
done

```

theorem $Goedel_II: \neg\ \{\} \vdash Fls \implies \neg\ \{\} \vdash neg\ (Pfp\ \langle Fls \rangle)$

by $(rule\ g2.goedel_second[unfolded\ consistent_def\ PP_def\ Pfp_subst\ subst.simps\ simp_thms\ if_True])$

lemma *ground_fm_PrjP[simp]*:
 $\text{ground_fm } (\text{PrjP } s \ k \ t) \longleftrightarrow \text{ground } s \wedge \text{ground } k \wedge \text{ground } t$
by (*auto simp add: ground_aux_def ground_fm_aux_def supp_conv_fresh*)

lemma *Fvars_HPPair[simp]*: $Fvars \ (\text{HPair } t \ u) = Fvars \ t \cup Fvars \ u$
unfolding *Fvars_def*
by *auto*

lemma *ground_HPPair[simp]*: $\text{ground } (\text{HPair } t \ u) \longleftrightarrow \text{ground } t \wedge \text{ground } u$
unfolding *ground_Fvars*
by *auto*

interpretation *dwdf: Deduct2_with_False_Disj* **where**

var = UNIV :: name set
and *trm = UNIV :: tm set*
and *fmla = UNIV :: fm set*
and *num = {t. ground t}*
and *Var = Var*
and *FvarsT = Fvars*
and *substT = $\lambda t \ u \ x. \text{subst } x \ u \ t$*
and *Fvars = Fvars*
and *subst = $\lambda A \ u \ x. \text{subst_fm } A \ x \ u$*
and *eql = (EQ)*
and *cnj = (AND)*
and *dsj = (OR)*
and *imp = (IMP)*
and *all = All*
and *exi = Ex*
and *fls = Fls*
and *prv = (\vdash) {}*
and *bprv = (\vdash) {}*
apply *unfold_locales*
subgoal **by** *simp*
subgoal **by** *simp*
subgoal **by** *simp*
subgoal **by** (*auto intro: Disj_I1*)
subgoal **by** (*auto intro: Disj_I2*)
subgoal **by** (*auto intro: ContraAssume*)
subgoal **by** *simp*
done

interpretation *Minimal_Truth_Soundness* **where**

var = UNIV :: name set
and *trm = UNIV :: tm set*
and *fmla = UNIV :: fm set*
and *num = {t. ground t}*
and *Var = Var*
and *FvarsT = Fvars*
and *substT = $\lambda t \ u \ x. \text{subst } x \ u \ t$*
and *Fvars = Fvars*
and *subst = $\lambda A \ u \ x. \text{subst_fm } A \ x \ u$*
and *eql = (EQ)*
and *cnj = (AND)*
and *dsj = (OR)*
and *imp = (IMP)*
and *all = All*
and *exi = Ex*

```

and fls = Fls
and prv = ( $\vdash$ ) {}
and isTrue = eval_fm e0
apply unfold_locales
subgoal by (auto simp: Fls_def)
subgoal by simp
subgoal by (auto simp only: ex_eval_fm_iff_exists_tm eval_fm.simps(4) subst_fm.simps)
subgoal by (auto simp only: ex_eval_fm_iff_exists_tm)
subgoal by (simp add: neg_def)
subgoal by (auto dest: hfthm_sound)
done

```

```

lemma neg_Neg:
  {}  $\vdash$  neg  $\varphi$  IFF Neg  $\varphi$ 
unfolding neg_def
by (auto simp: Fls_def intro: ContraAssume)

```

```

lemma ground_aux_mono:  $A \subseteq B \implies$  ground_aux t A  $\implies$  ground_aux t B
unfolding ground_aux_def by auto

```

interpretation g1: Goedel_Form_Minimal_Truth_Soundness_HBL1iff_prv_Cmpl_Pf_Classical **where**

```

  var = UNIV :: name set
and trm = UNIV :: tm set
and fmla = UNIV :: fm set
and num = {t. ground t}
and Var = Var
and FvarsT = Fvars
and substT =  $\lambda$ t u x. subst x u t
and Fvars = Fvars
and subst =  $\lambda$ A u x. subst_fm A x u
and eql = (EQ)
and cnj = (AND)
and dsj = (OR)
and imp = (IMP)
and all = All
and exi = Ex
and fls = Fls
and prv = ( $\vdash$ ) {}
and bprv = ( $\vdash$ ) {}
and enc = quot
and S = KRP (quot (Var xx)) (Var xx) (Var yy)
and P = PfP (Var xx)
and isTrue = eval_fm e0
and Pf = Ex xx' (Ex yy' (Var yy EQ HPair (Var xx') (Var yy') AND PrfP (Var xx') (Var yy') (Var
xx)))
apply unfold_locales
subgoal by simp
subgoal unfolding Fvars_def by (auto simp: fresh_at_base(2))
subgoal for  $\varphi$ 
  unfolding Let_def
  supply PfP.simps[simp del]
  apply (subst psubst_eq_rawpsubst2) apply (simp_all add: PfP.simps[of yy' xx' quot  $\varphi$ , simplified])
  apply (auto simp: eqv_def)
  apply (rule Ex_I[of _ _ _ HPair (Var xx') (Var yy')])
  apply (subst subst_fm_Ex_with_renaming[of xx _ xx' yy]; (auto simp: Conj_eqvt))
  apply (subst subst_fm_Ex_with_renaming[of zz _ yy' yy]; (auto simp: Conj_eqvt HPair_eqvt
PrfP_eqvt))

```

```

apply (rule Ex_I[of _ _ _ (Var xx')]; auto)
apply (rule Ex_I[of _ _ _ (Var yy')]; auto)
apply (rule Ex_I[of _ _ _ (Var yy')]; auto)
apply (rule Ex_I[of _ _ _ (Var xx')]; auto)
done
subgoal
  by (auto simp: PP_def proved_iff_proved_PfP[symmetric])
subgoal for n  $\varphi$ 
  unfolding Let_def
  apply (subst (1 2) psubst_eq_rawpsubst2)
    apply (simp_all add: ground_Fvars)
  apply (rule impI)
  apply (rule Sigma_fm_imp_thm)
    apply (auto simp: ground_Fvars[symmetric] elim: ground_aux_mono[OF empty_subsetI])
  apply (auto simp: ground_aux_def ground_fm_aux_def supp_conv_fresh fresh_at_base Fvars_def)
  done
subgoal for  $\varphi$ 
  apply (rule NegNeg_D)
  apply (auto simp: PP_def dest!: Iff_MP_same[OF neg_Neg] Iff_MP_same[OF Neg_cong[OF neg_Neg]])
  done
done

```

theorem *Goedel_I*: $\exists \varphi. \neg \{ \} \vdash \varphi \wedge \neg \{ \} \vdash \text{Neg } \varphi \wedge \text{eval_fm } e0 \varphi$

by (*meson Iff_MP2_same g1.recover_proofs.goedel_first_classic_strong[OF consistent] neg_Neg*)

The following interpretation is redundant, because *Goedel_Form_Minimal_Truth_Soundness_HBL1iff_prv_Compl* (interpreted above) is a sublocale of *Goedel_Form_Classical_HBL1_rev_prv_Minimal_Truth_Soundness_TIP*. However, the latter requires less infrastructure (no Pf formula).

The definition of *isTrue* prevents Isabelle from noticing that the locale has already been interpreted via the above *g1* interpretation of *Goedel_Form_Minimal_Truth_Soundness_HBL1iff_prv_Compl_Pf_Classical*.

definition *isTrue* **where**

isTrue = *eval_fm e0*

interpretation *g1'*: *Goedel_Form_Classical_HBL1_rev_prv_Minimal_Truth_Soundness_TIP* **where**

```

  var = UNIV :: name set
and trm = UNIV :: tm set
and fmla = UNIV :: fm set
and num = {t. ground t}
and Var = Var
and FvarsT = Fvars
and substT =  $\lambda t u x. \text{subst } x u t$ 
and Fvars = Fvars
and subst =  $\lambda A u x. \text{subst\_fm } A x u$ 
and eql = (EQ)
and cnj = (AND)
and dsj = (OR)
and imp = (IMP)
and all = All
and exi = Ex
and fls = Fls
and prv = ( $\vdash$ ) { $\}$ 
and bprv = ( $\vdash$ ) { $\}$ 
and enc = quot
and S = KRP (quot (Var xx)) (Var xx) (Var yy)
and P = PfP (Var xx)
and isTrue = isTrue
apply unfold_locales
unfolding isTrue_def

```

```

subgoal by (auto simp: Fls_def)
subgoal by simp
subgoal by (auto simp only: ex_eval_fm_iff_exists_tm eval_fm.simps(4) subst_fm.simps)
subgoal by (auto simp only: ex_eval_fm_iff_exists_tm)
subgoal by (simp add: neg_def)
subgoal by (auto dest: hfthm_sound)
subgoal by (auto simp: proved_iff_proved_PfP[symmetric] PP_def quot_fm_coding
  simp del: eval_fm_PfP
  dest!: Iff_MP_same[OF neg_Neg] Iff_MP_same[OF Neg_cong[OF neg_Neg]] NegNeg_D
  Sigma_fm_imp_thm[rotated 2])
done

```

```

theorem Goedel_I':  $\exists \varphi. \neg \{ \} \vdash \varphi \wedge \neg \{ \} \vdash \text{Neg } \varphi \wedge \text{isTrue } \varphi$ 
by (meson Iff_MP2_same g1'.goedel_first_classic_strong[OF consistent] neg_Neg)

```