

# Euler's Partition Theorem

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## Abstract

Euler's Partition Theorem states that the number of partitions with only distinct parts is equal to the number of partitions with only odd parts. The combinatorial proof follows John Harrison's pre-existing HOL Light formalization [1]. To understand the rough idea of the proof, I read the lecture notes of the MIT course 18.312 on Algebraic Combinatorics [2] by Gregg Musiker. This theorem is the 45th theorem of the Top 100 Theorems list.

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## 1 Euler's Partition Theorem

**theory** *Euler-Partition*

**imports**

*Main*

*Card-Number-Partitions.Number-Partition*

**begin**

### 1.1 Preliminaries

#### 1.1.1 Additions to Divides Theory

**lemma** *power-div-nat*:

**assumes**  $c \leq b$   
**assumes**  $a > 0$   
**shows**  $(a :: \text{nat}) \wedge b \text{ div } a \wedge c = a \wedge (b - c)$   
**by** (*metis* *assms* *nonzero-mult-div-cancel-right* *le-add-diff-inverse2* *less-not-refl2* *power-add* *power-not-zero*)

### 1.1.2 Additions to Groups-Big Theory

**lemma** *sum-div*:  
**assumes** *finite*  $A$   
**assumes**  $\bigwedge a. a \in A \implies (b :: 'b :: \text{euclidean-semiring}) \text{ dvd } f a$   
**shows**  $(\sum a \in A. f a) \text{ div } b = (\sum a \in A. (f a) \text{ div } b)$   
**using** *assms*  
**proof** (*induct*)  
**case** *insert* **from** *this* **show** *?case* **by** *auto* (*subst* *div-add*; *auto* *intro!*: *dvd-sum*)  
**qed** (*auto*)

**lemma** *sum-mod*:  
**assumes** *finite*  $A$   
**assumes**  $\bigwedge a. a \in A \implies f a \text{ mod } b = (0 :: 'b :: \text{unique-euclidean-semiring})$   
**shows**  $(\sum a \in A. f a) \text{ mod } b = 0$   
**using** *assms* **by** *induct* (*auto* *simp* *add*: *mod-add-eq* [*symmetric*])

### 1.1.3 Additions to Finite-Set Theory

**lemma** *finite-exponents*:  
*finite*  $\{i. 2 \wedge i \leq (n :: \text{nat})\}$   
**proof** –  
**have**  $\{i :: \text{nat}. 2 \wedge i \leq n\} \subseteq \{0..n\}$   
**using** *dual-order.trans* **by** *fastforce*  
**from** *finite-subset[OF this]* **show** *?thesis* **by** *simp*  
**qed**

## 1.2 Binary Encoding of Natural Numbers

**definition** *bitset* ::  $\text{nat} \Rightarrow \text{nat set}$   
**where**  
 $\text{bitset } n = \{i. \text{odd } (n \text{ div } (2 \wedge i))\}$

**lemma** *in-bitset-bound*:  
 $b \in \text{bitset } n \implies 2 \wedge b \leq n$   
**unfolding** *bitset-def* **using** *not-less* **by** *fastforce*

**lemma** *in-bitset-bound-weak*:  
 $b \in \text{bitset } n \implies b \leq n$   
**by** (*meson* *order.trans* *in-bitset-bound* *self-le-ge2-pow[OF order-refl]*)

**lemma** *finite-bitset*:  
*finite* (*bitset*  $n$ )  
**proof** –

**have**  $\text{bitset } n \subseteq \{..n\}$  **by** (*auto dest: in-bitset-bound-weak*)  
**from this show** *?thesis using finite-subset by auto*  
**qed**

**lemma** *bitset-0:*  
 $\text{bitset } 0 = \{\}$   
**unfolding** *bitset-def by auto*

**lemma** *bitset-2n:*  $\text{bitset } (2 * n) = \text{Suc } \text{' } (\text{bitset } n)$   
**proof** (*rule set-eqI*)  
**fix**  $x$   
**show**  $(x \in \text{bitset } (2 * n)) = (x \in \text{Suc } \text{' } \text{bitset } n)$   
**unfolding** *bitset-def by (cases x) auto*  
**qed**

**lemma** *bitset-Suc:*  
**assumes** *even n*  
**shows**  $\text{bitset } (n + 1) = \text{insert } 0 (\text{bitset } n)$   
**proof** (*rule set-eqI*)  
**fix**  $x$   
**from** *assms show*  $(x \in \text{bitset } (n + 1)) = (x \in \text{insert } 0 (\text{bitset } n))$   
**unfolding** *bitset-def by (cases x) (auto simp add: div-mult2-eq)*  
**qed**

**lemma** *bitset-2n1:*  
 $\text{bitset } (2 * n + 1) = \text{insert } 0 (\text{Suc } \text{' } (\text{bitset } n))$   
**by** (*subst bitset-Suc (auto simp add: bitset-2n)*)

**lemma** *sum-bitset:*  
 $(\sum_{i \in \text{bitset } n} 2^i) = n$   
**proof** (*induct rule: nat-bit-induct*)  
**case zero**  
**show** *?case by (auto simp add: bitset-0)*  
**next**  
**case (even n)**  
**from this show** *?case*  
**by** (*simp add: bitset-2n sum.reindex sum-distrib-left[symmetric]*)  
**next**  
**case (odd n)**  
**have**  $(\sum_{i \in \text{bitset } (2 * n + 1)} 2^i) = (\sum_{i \in \text{insert } 0 (\text{Suc } \text{' } \text{bitset } n)} 2^i)$   
**by** (*simp only: bitset-2n1*)  
**also have**  $\dots = 2^0 + (\sum_{i \in \text{Suc } \text{' } \text{bitset } n} 2^i)$   
**by** (*subst sum.insert (auto simp add: finite-bitset)*)  
**also have**  $\dots = 2 * n + 1$   
**using odd by** (*simp add: sum.reindex sum-distrib-left[symmetric]*)  
**finally show** *?case by simp*  
**qed**

**lemma** *binarysum-div:*

```

assumes finite B
shows  $(\sum i \in B. (2::nat) \wedge i) \text{ div } 2 \wedge j = (\sum i \in B. \text{if } i < j \text{ then } 0 \text{ else } 2 \wedge (i - j))$ 
(is - =  $(\sum i \in -. ?f i)$ )
proof -
  have split-B:  $B = \{i \in B. i < j\} \cup \{i \in B. j \leq i\}$  by auto
  have bound:  $(\sum i \mid i \in B \wedge i < j. (2::nat) \wedge i) < 2 \wedge j$ 
  proof (rule order.strict-trans1)
    show  $(\sum i \mid i \in B \wedge i < j. (2::nat) \wedge i) \leq (\sum i < j. 2 \wedge i)$  by (auto intro: sum-mono2)
    show  $\dots < 2 \wedge j$  using sum-power2 by (simp add: atLeast0LessThan)
  qed
  from this have zero:  $(\sum i \mid i \in B \wedge i < j. (2::nat) \wedge i) \text{ div } (2 \wedge j) = 0$  by (elim div-less)
  from assms have mod0:  $(\sum i \mid i \in B \wedge j \leq i. (2::nat) \wedge i) \text{ mod } 2 \wedge j = 0$ 
    by (auto intro!: sum-mod simp add: le-imp-power-dvd)
  from assms have  $(\sum i \in B. (2::nat) \wedge i) \text{ div } (2 \wedge j) = ((\sum i \mid i \in B \wedge i < j. 2 \wedge i) + (\sum i \mid i \in B \wedge j \leq i. 2 \wedge i)) \text{ div } 2 \wedge j$ 
    by (subst sum.union-disjoint[symmetric]) (auto simp add: split-B[symmetric])
  also have  $\dots = (\sum i \mid i \in B \wedge j \leq i. 2 \wedge i) \text{ div } 2 \wedge j$ 
    by (simp add: div-add1-eq zero mod0)
  also have  $\dots = (\sum i \mid i \in B \wedge j \leq i. 2 \wedge i \text{ div } 2 \wedge j)$ 
    using assms by (subst sum-div) (auto simp add: sum-div le-imp-power-dvd)
  also have  $\dots = (\sum i \mid i \in B \wedge j \leq i. 2 \wedge (i - j))$ 
    by (rule sum.cong[OF refl]) (auto simp add: power-div-nat)
  also have  $\dots = (\sum i \in B. ?f i)$ 
    using assms by (subst split-B; subst sum.union-disjoint) auto
  finally show ?thesis .
qed

lemma odd-iff:
  assumes finite B
  shows  $\text{odd } (\sum i \in B. \text{if } i < x \text{ then } (0::nat) \text{ else } 2 \wedge (i - x)) = (x \in B)$  (is odd  $(\sum i \in -. ?s i) = -$ )
proof -
  from assms have even:  $\text{even } (\sum i \in B - \{x\}. ?s i)$ 
    by (subst dvd-sum) auto
  show ?thesis
  proof
    assume odd  $(\sum i \in B. ?s i)$ 
    from this even show  $x \in B$  by (cases  $x \in B$ ) auto
  next
    assume  $x \in B$ 
    from assms this have  $(\sum i \in B. ?s i) = 1 + (\sum i \in B - \{x\}. ?s i)$ 
      by (auto simp add: sum.remove)
    from assms this even show  $\text{odd } (\sum i \in B. ?s i)$  by auto
  qed
qed

```

**lemma** *bitset-sum*:  
 assumes *finite B*  
 shows *bitset* ( $\sum_{i \in B}. 2^i$ ) = *B*  
 using *assms unfolding bitset-def* **by** (*simp add: binarysum-div odd-iff*)

### 1.3 Decomposition of a Number into a Power of Two and an Odd Number

**function** (*sequential*) *index* :: *nat*  $\Rightarrow$  *nat*  
**where**  
*index* 0 = 0  
 | *index* n = (*if odd n then 0 else Suc (index (n div 2))*)  
**by** (*pat-completeness*) *auto*

**termination**  
**proof**  
 show *wf* {(*x::nat*, *y*). *x* < *y*} **by** (*simp add: wf*)  
**next**  
 fix *n* show (*Suc n div 2*, *Suc n*)  $\in$  {(*x*, *y*). *x* < *y*} **by** *simp*  
**qed**

**function** (*sequential*) *oddp* :: *nat*  $\Rightarrow$  *nat*  
**where**  
*oddp* 0 = 0  
 | *oddp* n = (*if odd n then n else oddpart (n div 2)*)  
**by** *pat-completeness auto*

**termination**  
**proof**  
 show *wf* {(*x::nat*, *y*). *x* < *y*} **by** (*simp add: wf*)  
**next**  
 fix *n* show (*Suc n div 2*, *Suc n*)  $\in$  {(*x*, *y*). *x* < *y*} **by** *simp*  
**qed**

**lemma** *odd-oddp*:  
*odd (oddp n)*  $\longleftrightarrow$  *n*  $\neq$  0  
**by** (*induct n rule: index.induct*) *auto*

**lemma** *index-oddp-decomposition*:  
*n* =  $2^{\text{index } n} * \text{oddp } n$   
**proof** (*induct n rule: index.induct*)  
 case (2 *n*)  
 from *this* show *Suc n* =  $2^{\text{index } (\text{Suc } n)} * \text{oddp } (\text{Suc } n)$   
 by (*simp add: mult.assoc*)  
**qed** (*simp*)

**lemma** *oddp-leq*:  
*oddp n*  $\leq$  *n*  
**by** (*induct n rule: index.induct*) (*simp,metis div-le-dividend le-Suc-eq le-trans odd-*

*part.simps(2)*)

**lemma** *index-oddpart-unique*:  
 **assumes** *odd (m :: nat) odd m'*  
 **shows**  $(2 \wedge i * m = 2 \wedge i' * m') \longleftrightarrow (i = i' \wedge m = m')$   
**proof** (*induct i arbitrary: i'*)  
 **case** 0  
 **from** *assms show ?case by auto*  
**next**  
 **case** (*Suc - i'*)  
 **from** *assms this show ?case by (cases i') auto*  
**qed**

**lemma** *index-oddpart*:  
 **assumes** *odd m*  
 **shows**  $\text{index } (2 \wedge i * m) = i \text{ oddpart } (2 \wedge i * m) = m$   
**using** *index-oddpart-unique*[**where** *i=i and m=m and m'=oddpart (2 \wedge i \* m)*  
**and** *i'=index (2 \wedge i \* m)*]  
 *assms odd-oddpart index-oddpart-decomposition by force+*

## 1.4 Partitions With Only Distinct and Only Odd Parts

**definition** *odd-of-distinct* ::  $(\text{nat} \Rightarrow \text{nat}) \Rightarrow \text{nat} \Rightarrow \text{nat}$   
**where**  
  $\text{odd-of-distinct } p = (\lambda i. \text{if odd } i \text{ then } (\sum j \mid p (2 \wedge j * i) = 1. 2 \wedge j) \text{ else } 0)$

**definition** *distinct-of-odd* ::  $(\text{nat} \Rightarrow \text{nat}) \Rightarrow \text{nat} \Rightarrow \text{nat}$   
**where**  
  $\text{distinct-of-odd } p = (\lambda i. \text{if index } i \in \text{bitset } (p (\text{oddpart } i)) \text{ then } 1 \text{ else } 0)$

**lemma** *odd*:  
  $\text{odd-of-distinct } p \ i \neq 0 \implies \text{odd } i$   
**unfolding** *odd-of-distinct-def* **by** *auto*

**lemma** *distinct-distinct-of-odd*:  
  $\text{distinct-of-odd } p \ i \leq 1$   
**unfolding** *distinct-of-odd-def* **by** *auto*

**lemma** *odd-of-distinct*:  
 **assumes** *odd-of-distinct p i \neq 0*  
 **assumes**  $\bigwedge i. p \ i \neq 0 \implies i \leq n$   
 **shows**  $1 \leq i \wedge i \leq n$   
**proof**  
 **from** *assms(1) odd* **have** *odd i*  
 **by** *simp*  
 **then** **show**  $1 \leq i$   
 **by** (*auto elim: oddE*)  
**next**  
 **from** *assms(1)* **obtain** *j* **where**  $p (2 \wedge j * i) > 0$

by (auto simp add: odd-of-distinct-def split: if-splits) fastforce  
 with *assms*(2) have  $i \leq 2^j * i \wedge 2^j * i \leq n$   
 by *simp-all*  
 then show  $i \leq n$   
 by (rule order-trans)  
 qed

**lemma** *distinct-of-odd*:

assumes  $\bigwedge i. p\ i * i \leq n \wedge i. p\ i \neq 0 \implies \text{odd } i$

assumes *distinct-of-odd*  $p\ i \neq 0$

shows  $1 \leq i \wedge i \leq n$

**proof**

from *assms*(3) have *index*:  $i \in \text{bitset } (p\ (\text{oddpart } i))$

unfolding *distinct-of-odd-def* by (auto split: if-split-asm)

have  $i \neq 0$

**proof**

assume *zero*:  $i = 0$

from *assms*(2) have  $p\ 0 = 0$  by *auto*

from *index zero* this show *False* by (auto simp add: *bitset-0*)

qed

from this show  $1 \leq i$  by *auto*

from *assms*(1) have *leq-n*:  $p\ (\text{oddpart } i) * \text{oddpart } i \leq n$  by *auto*

from *index* have  $2^{\text{index } i} \leq p\ (\text{oddpart } i)$  by (rule *in-bitset-bound*)

from this *leq-n* show  $i \leq n$

by (*subst index-oddpart-decomposition*[of *i*]) (*meson dual-order.trans eq-imp-le mult-le-mono*)

qed

**lemma** *odd-distinct*:

assumes  $\bigwedge i. p\ i \neq 0 \implies \text{odd } i$

shows *odd-of-distinct* (*distinct-of-odd*  $p$ ) =  $p$

using *assms* unfolding *odd-of-distinct-def* *distinct-of-odd-def*

by (auto simp add: *fun-eq-iff index-oddpart sum-bitset*)

**lemma** *distinct-odd*:

assumes  $\bigwedge i. p\ i \neq 0 \implies 1 \leq i \wedge i \leq n \wedge i. p\ i \leq 1$

shows *distinct-of-odd* (*odd-of-distinct*  $p$ ) =  $p$

**proof** –

from *assms* have  $\{i. p\ i = 1\} \subseteq \{..n\}$  by *auto*

from this have *finite*: *finite*  $\{i. p\ i = 1\}$  by (*simp add: finite-subset*)

have  $\bigwedge x\ j. x > 0 \implies p\ (2^j * \text{oddpart } x) = 1 \implies$

*index*  $(2^j * \text{oddpart } x) \in \text{index } \{i. p\ i = 1 \wedge \text{oddpart } x = \text{oddpart } i\}$

by (rule *imageI*) (auto intro: *imageI simp add: index-oddpart odd-oddpart*)

from this have *eq*:  $\bigwedge x. x > 0 \implies \{j. p\ (2^j * \text{oddpart } x) = 1\} = \text{index } \{i. p\ i = 1 \wedge \text{oddpart } x = \text{oddpart } i\}$

by (auto simp add: *index-oddpart odd-oddpart index-oddpart-decomposition*[*symmetric*])

from *finite* have *all-finite*:  $\bigwedge x. x > 0 \implies \text{finite } \{j. p\ (2^j * \text{oddpart } x) = 1\}$

unfolding *eq* by *auto*

show *?thesis*

```

proof
  fix  $x$ 
  from  $assms(1)$  have  $p0: p\ 0 = 0$  by  $auto$ 
  show  $distinct\text{-of}\text{-odd}\ (odd\text{-of}\text{-distinct}\ p)\ x = p\ x$ 
  proof ( $cases\ x > 0$ )
    case  $False$ 
    from  $this\ p0$  show  $?thesis$ 
    unfolding  $odd\text{-of}\text{-distinct}\text{-def}\ distinct\text{-of}\text{-odd}\text{-def}$ 
    by ( $auto\ simp\ add: odd\text{-oddpart}\ bitset\text{-}0$ )
  next
  case  $True$ 
  from  $p0\ assms(2)[of\ x]\ all\text{-finite}[OF\ True]$  show  $?thesis$ 
  unfolding  $odd\text{-of}\text{-distinct}\text{-def}\ distinct\text{-of}\text{-odd}\text{-def}$ 
  by ( $auto\ simp\ add: odd\text{-oddpart}\ bitset\text{-}0\ bitset\text{-sum}\ index\text{-oddpart}\text{-decomposition}[symmetric]$ )
  qed
qed
qed

```

**lemma**  $sum\text{-distinct}\text{-of}\text{-odd}$ :

```

assumes  $\bigwedge i. p\ i \neq 0 \implies 1 \leq i \wedge i \leq n$ 
assumes  $\bigwedge i. p\ i * i \leq n$ 
assumes  $\bigwedge i. p\ i \neq 0 \implies odd\ i$ 
shows  $(\sum_{i \leq n}. distinct\text{-of}\text{-odd}\ p\ i * i) = (\sum_{i \leq n}. p\ i * i)$ 
proof -
  {
    fix  $m$ 
    assume  $odd: odd\ (m :: nat)$ 
    have  $finite: finite\ \{k. 2^k * m \leq n \wedge k \in bitset\ (p\ m)\}$  by ( $simp\ add: finite\text{-bitset}$ )
    have  $(\sum i \mid \exists k. i = 2^k * m \wedge i \leq n. distinct\text{-of}\text{-odd}\ p\ i * i) =$ 
       $(\sum i \mid \exists k. i = 2^k * m \wedge i \leq n. if\ index\ i \in bitset\ (p\ (oddpart\ i))\ then\ i$ 
       $else\ 0)$ 
    unfolding  $distinct\text{-of}\text{-odd}\text{-def}$  by ( $auto\ intro: sum.cong$ )
    also have  $\dots = (\sum i \mid \exists k. i = 2^k * m \wedge k \in bitset\ (p\ m) \wedge i \leq n. i)$ 
    using  $odd$  by ( $intro\ sum.mono\text{-neutral}\text{-cong}\text{-right}$ ) ( $auto\ simp\ add: index\text{-oddpart}$ )
    also have  $\dots = (\sum k \mid 2^k * m \leq n \wedge k \in bitset\ (p\ m). 2^k * m)$ 
    using  $odd$  by ( $auto\ intro!: sum.reindex\text{-cong}[OF\ -\ refl]\ inj\text{-on}I$ )
    also have  $\dots = (\sum_{k \in bitset\ (p\ m)}. 2^k * m)$ 
    using  $assms(2)[of\ m]\ finite\ dual\text{-order.trans}\ in\text{-bitset}\text{-bound}$ 
    by ( $fastforce\ intro!: sum.mono\text{-neutral}\text{-cong}\text{-right}$ )
    also have  $\dots = (\sum_{k \in bitset\ (p\ m)}. 2^k) * m$ 
    by ( $subst\ sum\text{-distrib}\text{-right}$ )  $auto$ 
    also have  $\dots = p\ m * m$ 
    by ( $auto\ simp\ add: sum\text{-bitset}$ )
    finally have  $(\sum i \mid \exists k. i = 2^k * m \wedge i \leq n. distinct\text{-of}\text{-odd}\ p\ i * i) = p\ m$ 
     $*\ m.$ 
  } note  $inner\text{-eq} = this$ 

```

**have** *set-eq*:  $\{i. 1 \leq i \wedge i \leq n\} = \bigcup ((\lambda m. \{i. \exists k. i = (2 \wedge k) * m \wedge i \leq n\}) \text{ ‘ } \{m. m \leq n \wedge \text{odd } m\})$   
**proof** –  
{  
  **fix** *x*  
  **assume**  $1 \leq x \wedge x \leq n$   
  **from** *this oddpart-leq*[of *x*] **have**  $\text{oddpart } x \leq n \wedge \text{odd } (\text{oddpart } x) \wedge (\exists k. 2 \wedge k \text{ index } x * \text{oddpart } x = 2 \wedge k * \text{oddpart } x)$   
  **by** (*auto simp add: odd-oddpart*)  
  **from** *this* **have**  $\exists m \leq n. \text{odd } m \wedge (\exists k. x = 2 \wedge k * m)$   
  **by** (*auto simp add: index-oddpart-decomposition[symmetric]*)  
}  
**from** *this* **show** *?thesis* **by** (*auto simp add: Suc-leI odd-pos*)  
**qed**  
**let** *?S* =  $(\lambda m. \{i. \exists k. i = 2 \wedge k * m \wedge i \leq n\}) \text{ ‘ } \{m. m \leq n \wedge \text{odd } m\}$   
**have** *no-overlap*:  $\forall A \in ?S. \forall B \in ?S. A \neq B \longrightarrow A \cap B = \{\}$   
**by** (*auto simp add: index-oddpart-unique*)  
**have** *inj*: *inj-on*  $(\lambda m. \{i. (\exists k. i = 2 \wedge k * m) \wedge i \leq n\}) \{m. m \leq n \wedge \text{odd } m\}$   
**unfolding** *inj-on-def* **by** *auto* (*force simp add: index-oddpart-unique*)  
**have** *reindex*:  $\bigwedge F. (\sum i \mid 1 \leq i \wedge i \leq n. F i) = (\sum m \mid m \leq n \wedge \text{odd } m. (\sum i \mid \exists k. i = 2 \wedge k * m \wedge i \leq n. F i))$   
**unfolding** *set-eq* **by** (*subst sum.Union-disjoint*) (*auto simp add: no-overlap intro: sum.reindex-cong[OF inj]*)  
**have**  $(\sum i \leq n. \text{distinct-of-odd } p i * i) = (\sum i \mid 1 \leq i \wedge i \leq n. \text{distinct-of-odd } p i * i)$   
**by** (*auto intro: sum.mono-neutral-right*)  
**also** **have**  $\dots = (\sum m \mid m \leq n \wedge \text{odd } m. \sum i \mid \exists k. i = 2 \wedge k * m \wedge i \leq n. \text{distinct-of-odd } p i * i)$   
**by** (*simp only: reindex*)  
**also** **have**  $\dots = (\sum i \mid i \leq n \wedge \text{odd } i. p i * i)$   
**by** (*rule sum.cong[OF refl]; subst inner-eq*) *auto*  
**also** **have**  $\dots = (\sum i \leq n. p i * i)$   
**using** *assms*(3) **by** (*auto intro: sum.mono-neutral-left*)  
**finally** **show** *?thesis* .  
**qed**

**lemma** *leq-n*:

**assumes**  $\forall i. 0 < p i \longrightarrow 1 \leq i \wedge i \leq (n::\text{nat})$   
**assumes**  $(\sum i \leq n. p i * i) = n$   
**shows**  $p i * i \leq n$   
**proof** (*rule ccontr*)  
**assume**  $\neg p i * i \leq n$   
**from** *this* **have** *gr-n*:  $p i * i > n$  **by** *auto*  
**from** *this* *assms*(1) **have**  $1 \leq i \wedge i \leq n$  **by** *force*  
**from** *this* **have**  $(\sum j \leq n. p j * j) = p i * i + (\sum j \mid j \leq n \wedge j \neq i. p j * j)$   
**by** (*subst sum.insert[symmetric]*) (*auto intro: sum.cong simp del: sum.insert*)  
**from** *this* *gr-n* *assms*(2) **show** *False* **by** *simp*  
**qed**

**lemma** *distinct-of-odd-in-distinct-partitions*:  
**assumes**  $p \in \{p. p \text{ partitions } n \wedge (\forall i. p \ i \neq 0 \longrightarrow \text{odd } i)\}$   
**shows**  $\text{distinct-of-odd } p \in \{p. p \text{ partitions } n \wedge (\forall i. p \ i \leq 1)\}$   
**proof**  
**have** *distinct-of-odd p partitions n*  
**proof** (*rule partitionsI*)  
**fix**  $i$  **assume** *distinct-of-odd p i  $\neq 0$*   
**from** *this assms* **show**  $1 \leq i \wedge i \leq n$   
**unfolding** *partitions-def*  
**by** (*rule-tac distinct-of-odd*) (*auto simp add: leq-n*)  
**next**  
**from** *assms* **show**  $(\sum_{i \leq n}. \text{distinct-of-odd } p \ i * i) = n$   
**by** (*subst sum-distinct-of-odd*) (*auto simp add: distinct-distinct-of-odd leq-n elim: partitionsE*)  
**qed**  
**moreover** **have**  $\forall i. \text{distinct-of-odd } p \ i \leq 1$   
**by** (*intro allI distinct-distinct-of-odd*)  
**ultimately** **show**  $\text{distinct-of-odd } p \text{ partitions } n \wedge (\forall i. \text{distinct-of-odd } p \ i \leq 1)$   
**by** *simp*  
**qed**

**lemma** *odd-of-distinct-in-odd-partitions*:  
**assumes**  $p \in \{p. p \text{ partitions } n \wedge (\forall i. p \ i \leq 1)\}$   
**shows**  $\text{odd-of-distinct } p \in \{p. p \text{ partitions } n \wedge (\forall i. p \ i \neq 0 \longrightarrow \text{odd } i)\}$   
**proof**  
**from** *assms* **have** *distinct:  $\bigwedge i. p \ i = 0 \vee p \ i = 1$*   
**using** *le-imp-less-Suc less-Suc-eq-0-disj* **by** *fastforce*  
**from** *assms* **have** *set-eq:  $\{x. p \ x = 1\} = \{x \in \{..n\}. p \ x = 1\}$*   
**unfolding** *partitions-def* **by** *auto*  
**from** *assms* **have** *sum:  $(\sum_{i \leq n}. p \ i * i) = n$*   
**unfolding** *partitions-def* **by** *auto*  
**{**  
**fix**  $i$   
**assume**  $i: \text{odd } (i :: \text{nat})$   
**have**  $\exists: \text{inj-on index } \{x. p \ x = 1 \wedge \text{oddpart } x = i\}$   
**unfolding** *inj-on-def* **by** *auto* (*metis index-oddpart-decomposition*)  
**{**  
**fix**  $j$  **assume**  $p \ (2 \wedge^j * i) = 1$   
**from** *this i* **have**  $j \in \text{index } \{x. p \ x = 1 \wedge \text{oddpart } x = i\}$   
**by** (*auto simp add: index-oddpart(1, 2) intro!: image-eqI* [**where**  $x=2 \wedge^j * i$ ])  
**}**  
**from**  $i$  *this* **have**  $\{j. p \ (2 \wedge^j * i) = 1\} = \text{index } \{x. p \ x = 1 \wedge \text{oddpart } x = i\}$   
**by** (*auto simp add: index-oddpart-decomposition[symmetric]*)  
**from**  $\exists$  *this* **have**  $(\sum j \mid p \ (2 \wedge^j * i) = 1. 2 \wedge^j) * i = (\sum x \mid p \ x = 1 \wedge \text{oddpart } x = i. 2 \wedge^{\text{index } x} * i)$   
**by** (*auto intro: sum.reindex-cong* [**where**  $l = \text{index}$ ])  
**also** **have**  $\dots = (\sum x \mid p \ x = 1 \wedge \text{oddpart } x = i. 2 \wedge^{\text{index } x} * \text{oddpart } x)$   
**by** (*auto simp add: sum-distrib-right*)

**also have**  $\dots = (\sum x \mid p\ x = 1 \wedge \text{oddpart } x = i.\ x)$   
**by** (*simp only: index-oddpart-decomposition[symmetric]*)  
**also have**  $\dots \leq (\sum x \mid p\ x = 1.\ x)$   
**using** *set-eq* **by** (*intro sum-mono2*) *auto*  
**also have**  $\dots = (\sum x \leq n.\ p\ x * x)$   
**using** *distinct* **by** (*subst set-eq*) (*force intro!: sum.mono-neutral-cong-left*)  
**also have**  $\dots = n$  **using** *sum* .  
**finally have**  $(\sum j \mid p\ (2^j * i) = 1.\ 2^j * i \leq n)$  .  
**}**  
**from this have** *less-n:  $\bigwedge i.\ \text{odd-of-distinct } p\ i * i \leq n$*   
**unfolding** *odd-of-distinct-def* **by** *auto*  
**have** *odd-of-distinct p partitions n*  
**proof** (*rule partitionsI*)  
**fix** *i* **assume** *odd-of-distinct p i  $\neq 0$*   
**from this assms show**  $1 \leq i \wedge i \leq n$   
**by** (*elim CollectE conjE partitionsE odd-of-distinct*) *auto*  
**next**  
**have**  $(\sum i \leq n.\ \text{odd-of-distinct } p\ i * i) = (\sum i \leq n.\ \text{distinct-of-odd } (\text{odd-of-distinct } p)\ i * i)$   
**using** *assms less-n* **by** (*subst sum-distinct-of-odd*) (*auto elim!: partitionsE odd-of-distinct simp only: odd*)  
**also have**  $\dots = (\sum i \leq n.\ p\ i * i)$  **using** *assms*  
**by** (*auto elim!: partitionsE simp only:*) (*subst distinct-odd, auto*)  
**also with assms have**  $\dots = n$  **by** (*auto elim: partitionsE*)  
**finally show**  $(\sum i \leq n.\ \text{odd-of-distinct } p\ i * i) = n$  .  
**qed**  
**moreover have**  $\forall i.\ \text{odd-of-distinct } p\ i \neq 0 \longrightarrow \text{odd } i$   
**by** (*intro allI impI odd*)  
**ultimately show** *odd-of-distinct p partitions n  $\wedge (\forall i.\ \text{odd-of-distinct } p\ i \neq 0 \longrightarrow \text{odd } i)$*  **by** *simp*  
**qed**

## 1.5 Euler's Partition Theorem

**theorem** *Euler-partition-theorem:*

$\text{card } \{p.\ p\ \text{partitions } n \wedge (\forall i.\ p\ i \leq 1)\} = \text{card } \{p.\ p\ \text{partitions } n \wedge (\forall i.\ p\ i \neq 0 \longrightarrow \text{odd } i)\}$

(*is card ?distinct-partitions = card ?odd-partitions*)

**proof** (*rule card-bij-eq*)

**from** *odd-of-distinct-in-odd-partitions* **show**

*odd-of-distinct ' ?distinct-partitions  $\subseteq$  ?odd-partitions* **by** *auto*

**moreover from** *distinct-of-odd-in-distinct-partitions* **show**

*distinct-of-odd ' ?odd-partitions  $\subseteq$  ?distinct-partitions* **by** *auto*

**moreover have**  $\forall p \in ?\text{distinct-partitions}.\ \text{distinct-of-odd } (\text{odd-of-distinct } p) = p$

**by** *auto* (*subst distinct-odd; auto simp add: partitions-def*)

**moreover have**  $\forall p \in ?\text{odd-partitions}.\ \text{odd-of-distinct } (\text{distinct-of-odd } p) = p$

**by** *auto* (*subst odd-distinct; auto simp add: partitions-def*)

**ultimately show** *inj-on odd-of-distinct ?distinct-partitions*

*inj-on distinct-of-odd ?odd-partitions*

```
by (intro bij-betw-imp-inj-on bij-betw-byWitness; auto)+
show finite ?distinct-partitions finite ?odd-partitions
by (simp add: finite-partitions)+
qed

end
```

## References

- [1] J. Harrison. Euler's partition theorem and other elementary partition theorems. <https://github.com/jrh13/hol-light/blob/master/100/euler.ml>.
- [2] G. Musiker. Course 18.312: Algebraic combinatorics, 2009. [http://ocw.mit.edu/courses/mathematics/18-312-algebraic-combinatorics-spring-2009/readings-and-lecture-notes/MIT18\\_312S09\\_lec10\\_Patitio.pdf](http://ocw.mit.edu/courses/mathematics/18-312-algebraic-combinatorics-spring-2009/readings-and-lecture-notes/MIT18_312S09_lec10_Patitio.pdf).