Dynamic Tables

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Abstract

This article formalizes the amortized analysis of dynamic tables parameterized with their minimal and maximal load factors and the expansion and contraction factors.

A full description is found in a companion paper [1].

theory *Tables-real*
imports *Amortized-Complexity.Amortized-Framework0*
begin

fun Ψ :: *bool* $\Rightarrow$ *real* $\Rightarrow$ *real* $\Rightarrow$ *real* $\Rightarrow$ *real* $\Rightarrow$ *nat* $\Rightarrow$ *real* **where**
$\Psi\ b\ i\ d\ x_1\ x_2\ n = (\mathit{if}\ n \geq x_2\ \mathit{then}\ i{*}(n - x_2)\ \mathit{else}$
$\mathit{if}\ n \leq x_1 \wedge b\ \mathit{then}\ d{*}(x_1 - n)\ \mathit{else}\ 0)$

declare *of-nat-Suc*[*simp*] *of-nat-diff* [*simp*]

An automatic proof:

lemma *Psi-diff-Ins*:
$0 < i \Longrightarrow 0 < d \Longrightarrow \Psi\ b\ i\ d\ x_1\ x_2\ (\mathit{Suc}\ n) - \Psi\ b\ i\ d\ x_1\ x_2\ n \leq i$
⟨*proof*⟩

lemma assumes [*arith*]: $0 < i\ 0 \leq d$
shows $\Psi\ b\ i\ d\ x_1\ x_2\ (n{+}1) - \Psi\ b\ i\ d\ x_1\ x_2\ n \leq i$ (**is** $?D \leq$ -)
⟨*proof*⟩

lemma *Psi-diff-Del*: **assumes** [*arith*]: $0 < i\ 0 \leq d\ n{\neq}0$ **and** $x_1 \leq x_2$
shows $\Psi\ b\ i\ d\ x_1\ x_2\ (n{-}\mathit{Suc}\ 0) - \Psi\ b\ i\ d\ x_1\ x_2\ (n) \leq d$ (**is** $?D \leq$ -)
⟨*proof*⟩

locale *Table0* =
fixes *f1 f2 f1' f2' e c* :: *real*
assumes *e1*[*arith*]: $e > 1$
assumes *c1*[*arith*]: $c > 1$
assumes *f1*[*arith*]: $f1 > 0$
assumes *f1cf2*: $f1{*}c < f2$
assumes *f1f2e*: $f1 < f2/e$

assumes *f1′-def*: $f1' = min\ (f1{*}c)\ (f2/e)$
assumes *f2′-def*: $f2' = max\ (f1{*}c)\ (f2/e)$
begin

lemma *f2*[*arith*]: $0 < f2$
⟨*proof*⟩

lemma *f2′*[*arith*]: $0 < f2'$
⟨*proof*⟩

lemma *f2′-less-f2*: $f2' < f2$
⟨*proof*⟩

lemma *f1-less-f1′*: $f1 < f1'$
⟨*proof*⟩

lemma *f1′-gr0*[*arith*]: $f1' > 0$
⟨*proof*⟩

lemma *f1′-le-f2′*: $f1' \leq f2'$
⟨*proof*⟩

lemma *f1′c-le-f1*: $f1'/c \leq f1$
⟨*proof*⟩

lemma *f2-le-f2′e*: $f2 \leq f2'{*}e$
⟨*proof*⟩

lemma *f1f2′c*: $f1 \leq f2'/c$
⟨*proof*⟩

lemma *f1′ef2*: $f1' * e \leq f2$
⟨*proof*⟩

end

locale *Table* = *Table0* +
fixes *l0* :: *real*
assumes *l0f2e*: $l0 \geq 1/(f2 * (e-1))$
assumes *l0f1c*: $l0 \geq 1/(f1 * (c-1))$
assumes *f2f2′*: $l0 \geq 1/(f2 - f2')$
assumes *f1′f1*: $l0 \geq 1/((f1' - f1){*}c)$
begin

definition $ai = f2/(f2-f2')$
definition $ad = f1/(f1'-f1)$

lemma *aigr0*[*arith*]: $ai > 1$
⟨*proof*⟩

lemma *adgr0*[*arith*]: $ad > 0$
⟨*proof*⟩

lemma *l0-gr0*[*arith*]: $l0 > 0$
⟨*proof*⟩

lemma *f1-l0*: **assumes** $l0 \leq l/c$ **shows** $f1*(l/c) \leq f1*l - 1$
⟨*proof*⟩

fun *nxt* :: $op_{tb} \Rightarrow nat*real \Rightarrow nat*real$ **where**
nxt Ins (*n*,*l*) =
 ($n+1$, *if* $n+1 \leq f2*l$ *then l else* $e*l$) |
nxt Del (*n*,*l*) =
 ($n-1$, *if* $f1*l \leq real(n-1)$ *then l else if* $l0 \leq l/c$ *then* l/c *else l*)

fun *T* :: $op_{tb} \Rightarrow nat*real \Rightarrow real$ **where**
T Ins (*n*,*l*) = (*if* $n+1 \leq f2*l$ *then 1 else* $n+1$) |
T Del (*n*,*l*) = (*if* $f1*l \leq real(n-1)$ *then 1 else if* $l0 \leq l/c$ *then n else 1*)

fun Φ :: $nat * real \Rightarrow real$ **where**
Φ (*n*,*l*) = (*if* $n \geq f2'*l$ *then* $ai*(n - f2'*l)$ *else*
  if $n \leq f1'*l \wedge l0 \leq l/c$ *then* $ad*(f1'*l - n)$ *else 0*)

lemma *Phi-Psi*: Φ (*n*,*l*) = Ψ ($l0 \leq l/c$) *ai ad* ($f1'*l$) ($f2'*l$) *n*
⟨*proof*⟩

fun *invar* **where**
invar(*n*,*l*) = ($l \geq l0 \wedge (l/c \geq l0 \longrightarrow f1*l \leq n) \wedge n \leq f2*l$)

abbreviation $U \equiv \lambda f$ -. *case f of Ins* ⇒ $ai+1$ | *Del* ⇒ $ad+1$

interpretation *tb*: *Amortized*
 where *init* = (*0*,*l0*) **and** *nxt* = *nxt*
 and *inv* = *invar*
 and *T* = *T* **and** Φ = Φ
 and *U* = *U*
⟨*proof*⟩

end

locale *Optimal* =
fixes *f2 c e* :: *real* **and** *l0* :: *nat*
assumes *e1*[*arith*]: $e > 1$
assumes *c1*[*arith*]: $c > 1$
assumes [*arith*]: $f2 > 0$
assumes *l0*: $(e*c)/(f2*(min\ e\ c - 1)) \leq l0$
begin

lemma *l0e*: $(e*c)/(f2*(e-1)) \le l0$
⟨*proof*⟩

lemma *l0c*: $(e*c)/(f2*(c-1)) \le l0$
⟨*proof*⟩

interpretation *Table*
where $f1=f2/(e*c)$ **and** $f2=f2$ **and** $e=e$ **and** $c=c$ **and** $f1'=f2/e$ **and** $f2'=f2/e$ **and** $l0=l0$
⟨*proof*⟩

lemma $ai = e/(e-1)$
⟨*proof*⟩

lemma $ad = 1/(c-1)$
⟨*proof*⟩

end

interpretation *I1*: *Optimal* **where** $e=2$ **and** $c=2$ **and** $f2=1$ **and** $l0=4$
⟨*proof*⟩

interpretation *I2*: *Optimal* **where** $e=2$ **and** $c=2$ **and** $f2=3/4$ **and** $l0=6$
⟨*proof*⟩

interpretation *I3*: *Optimal* **where** $e=2$ **and** $c=2$ **and** $f2=0.8$ **and** $l0=5$
⟨*proof*⟩

interpretation *I4*: *Optimal* **where** $e=3$ **and** $c=3$ **and** $f2=0.9$ **and** $l0=5$
⟨*proof*⟩

interpretation *I5*: *Optimal* **where** $e=4$ **and** $c=4$ **and** $f2=1$ **and** $l0=6$
⟨*proof*⟩

interpretation *I6*: *Optimal* **where** $e=2.5$ **and** $c=2.5$ **and** $f2=1$ **and** $l0=5$
⟨*proof*⟩

interpretation *I7*: *Optimal* **where** $f2=1$ **and** $c=3/2$ **and** $e=2$ **and** $l0=6$
⟨*proof*⟩

interpretation *I8*: *Optimal* **where** $f2=1$ **and** $e=3/2$ **and** $c=2$ **and** $l0=6$
⟨*proof*⟩

end
theory *Tables-nat*
imports *Tables-real*
begin

declare *le-of-int-ceiling*[*simp*]

locale *TableInv* = *Table0 f1 f2 f1′ f2′ e c* **for** *f1 f2 f1′ f2′ e c* :: *real* +
fixes *l0* :: *nat*
assumes *l0f2e*: $l0 \geq 1/(f2 * (e-1))$
assumes *l0f1c*: $l0 \geq 1/(f1 * (c-1))$

assumes *l0f2f1e*: $l0 \geq f1/(f2 - f1*e)$
assumes *l0f2f1c*: $l0 \geq f2/(f2 - f1*c)$
begin

lemma *l0-gr0*[*arith*]: $l0 > 0$
⟨*proof*⟩

lemma *f1-l0*: **assumes** $l0 \leq l/c$ **shows** $f1*(l/c) \leq f1*l - 1$
⟨*proof*⟩

fun *nxt* :: $op_{tb} \Rightarrow nat*nat \Rightarrow nat*nat$ **where**
nxt Ins (*n*,*l*) =
 ($n+1$, *if* $n+1 \leq f2*l$ *then l else nat*$\lceil e*l \rceil$) |
nxt Del (*n*,*l*) =
 ($n-1$, *if* $f1*l \leq real(n-1)$ *then l else if* $l0 \leq \lfloor l/c \rfloor$ *then nat*$\lfloor l/c \rfloor$ *else l*)

fun *T* :: $op_{tb} \Rightarrow nat*nat \Rightarrow real$ **where**
T Ins (*n*,*l*) = (*if* $n+1 \leq f2*l$ *then 1 else* $n+1$) |
T Del (*n*,*l*) = (*if* $f1*l \leq real(n-1)$ *then 1 else if* $l0 \leq \lfloor l/c \rfloor$ *then n else 1*)

fun *invar* :: $nat * nat \Rightarrow bool$ **where**
$invar(n,l) = (l \geq l0 \land (\lfloor l/c \rfloor \geq l0 \longrightarrow f1*l \leq n) \land n \leq f2*l)$

lemma *invar-init*: *invar* (*0*,*l0*)
⟨*proof*⟩

lemma *invar-pres*: **assumes** *invar s* **shows** *invar*(*nxt f s*)
⟨*proof*⟩

end

locale *Table1* = *TableInv* +
assumes *f2f2′*: $l0 \geq 1/(f2 - f2')$
assumes *f1′f1*: $l0 \geq 1/((f1' - f1)*c)$
begin

definition $ai = f2/(f2-f2')$
definition $ad = f1/(f1'-f1)$

lemma *aigr0*[*arith*]: $ai > 1$

⟨proof⟩

lemma *adgr0*[*arith*]: $ad > 0$
⟨proof⟩

lemma *f1'ad*[*arith*]: $f1'*ad > 0$
⟨proof⟩

lemma *f2'ai*[*arith*]: $f2'*ai > 0$
⟨proof⟩

fun Φ :: *nat* ∗ *nat* ⇒ *real* **where**
Φ $(n,l) = (if\ n \geq f2'*l\ then\ ai*(n - f2'*l)\ else$
  $if\ n \leq f1'*l \wedge l0 \leq \lfloor l/c \rfloor\ then\ ad*(f1'*l - n)\ else\ 0)$

lemma *Phi-Psi*: $\Phi\ (n,l) = \Psi\ (l0 \leq \lfloor l/c \rfloor)\ ai\ ad\ (f1'*l)\ (f2'*l)\ n$
⟨proof⟩

abbreviation $U \equiv \lambda f$ -. *case f of Ins* ⇒ $ai+1 + f1'*ad$ | *Del* ⇒ $ad+1 + f2'*ai$

interpretation *tb*: *Amortized*
  where *init* = (*0*,*l0*) **and** *nxt* = *nxt*
  and *inv* = *invar*
  and *T* = *T* **and** Φ = Φ
  and *U* = *U*
⟨proof⟩

end

locale *Table2-f1f2''* = *TableInv* +
fixes *f1'' f2''* :: *real*

locale *Table2* = *Table2-f1f2''* +
assumes *f2f2''*: $(f2 - f2'')*l0 \geq 1$
assumes *f1''f1*: $(f1'' - f1)*c*l0 \geq 1$

assumes *f1-less-f1''*: $f1 < f1''$
assumes *f1''-less-f1'*: $f1'' < f1'$
assumes *f2'-less-f2''*: $f2' < f2''$
assumes *f2''-less-f2*: $f2'' < f2$
assumes *f1''-f1'*: $l \geq real\ l0 \Longrightarrow f1'' * (l+1) \leq f1'*l$
assumes *f2'-f2''*: $l \geq real\ l0 \Longrightarrow f2' * l \leq f2'' * (l-1)$
begin

definition $ai = f2\ /\ (f2 - f2'')$
definition $ad = f1\ /\ (f1'' - f1)$

lemma *f1''-gr0*[*arith*]: $f1'' > 0$
⟨proof⟩

lemma *f2''-gr0*[*arith*]: $f2'' > 0$
⟨*proof*⟩

lemma *aigr0*[*arith*]: $ai > 0$
⟨*proof*⟩

lemma *adgr0*[*arith*]: $ad > 0$
⟨*proof*⟩

fun Φ :: *nat* ∗ *nat* ⇒ *real* **where**
$\Phi(n,l) = (\textit{if } n \geq f2''{*}l \textit{ then } ai{*}(n - f2''{*}l) \textit{ else}$
$\quad \textit{if } n \leq f1''{*}l \wedge l0 \leq \lfloor l/c \rfloor \textit{ then } ad{*}(f1''{*}l - n) \textit{ else } 0)$

lemma *Phi-Psi*: $\Phi\ (n,l) = \Psi\ (l0 \leq \lfloor l/c \rfloor)\ ai\ ad\ (f1''{*}l)\ (f2''{*}l)\ n$
⟨*proof*⟩

abbreviation $U \equiv \lambda f$ -. *case f of Ins* ⇒ *ai+1* | *Del* ⇒ *ad+1*

interpretation *tb*: *Amortized*
 where *init* = (*0*,*l0*) **and** *nxt* = *nxt*
 and *inv* = *invar*
 and *T* = *T* **and** Φ = Φ
 and *U* = *U*
⟨*proof*⟩

end

locale *Table3* = *Table2-f1f2''* +
assumes *f1''-def*: $f1'' = (f1'{::}real){*}l0/(l0+1)$
assumes *f2''-def*: $f2'' = (f2'{::}real){*}l0/(l0-1)$

assumes *l0-f2f2'*: $l0 \geq (f2+1)/(f2-f2')$
assumes *l0-f1f1'*: $l0 \geq (f1'{*}c+1)/((f1'-f1){*}c)$

assumes *l0-f1-f1'*: $l0 > f1/((f1'-f1))$
assumes *l0-f2-f2'*: $l0 > f2/(f2-f2')$
begin

lemma *l0-gr1*: $l0 > 1$
⟨*proof*⟩

lemma *f1''-less-f1'*: $f1'' < f1'$
⟨*proof*⟩

lemma *f1-less-f1''*: $f1 < f1''$

⟨*proof*⟩

lemma *f2′-less-f2″*: $f2' < f2''$
⟨*proof*⟩

lemma *f2″-less-f2*: $f2'' < f2$
⟨*proof*⟩

lemma *f2f2″*: $(f2 - f2'')*l0 \geq 1$
⟨*proof*⟩

lemma *f1″f1*: $(f1'' - f1)*c*l0 \geq 1$
⟨*proof*⟩

lemma *f1″-f1′*: **assumes** $l \geq real\ l0$ **shows** $f1''*(l+1) \leq f1' * l$
⟨*proof*⟩

lemma *f2′-f2″*: **assumes** $l \geq real\ l0$ **shows** $f2' * l \leq f2'' * (l-1)$
⟨*proof*⟩

sublocale *Table2*
⟨*proof*⟩

end

end

References

[1] T. Nipkow. Parameterized dynamic tables. http://www.in.tum.de/~nipkow/pubs/, 2015.