

# Positional Notation for Natural Numbers in an Arbitrary Base

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## Abstract

We demonstrate the existence and uniqueness of the base- $n$  representation of a natural number, where  $n$  is any natural number greater than 1. This comes up when trying to translate mathematical contest problems and solutions into Isabelle/HOL.

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**theory** *DigitsInBase*

**imports** *HOL-Computational-Algebra.Computational-Algebra HOL-Number-Theory.Number-Theory*

**begin**

## 1 Infinite sums

In this section, it is shown that an infinite series *of natural numbers* converges if and only if its terms are eventually zero. Additionally, the notion of a summation starting from an index other than zero is defined. A few obvious lemmas about these notions are established.

**definition** *eventually-zero* ::  $(\text{nat} \Rightarrow \text{-}) \Rightarrow \text{bool}$  **where**  
*eventually-zero* ( $D :: \text{nat} \Rightarrow \text{-}$ )  $\longleftrightarrow (\forall_{\infty} n. D\ n = 0)$

**lemma** *eventually-zero-char*:

**shows** *eventually-zero*  $D \longleftrightarrow (\exists s. \forall i \geq s. D\ i = 0)$   
 ⟨*proof*⟩

There's a lot of commonality between this setup and univariate polynomials, but drawing out the similarities and proving them is beyond the scope of the current version of this theory except for the following lemma.

**lemma** *eventually-zero-poly*:  
**shows** *eventually-zero*  $D \longleftrightarrow D = \text{poly.coeff } (Abs\text{-poly } D)$   
 ⟨*proof*⟩

**lemma** *eventually-zero-imp-summable* [*simp*]:  
**assumes** *eventually-zero*  $D$   
**shows** *summable*  $D$   
 ⟨*proof*⟩

**lemma** *summable-bounded*:  
**fixes**  $my\text{-seq} :: nat \Rightarrow nat$  **and**  $n :: nat$   
**assumes**  $\bigwedge i. i \geq n \longrightarrow my\text{-seq } i = 0$   
**shows** *summable*  $my\text{-seq}$   
 ⟨*proof*⟩

**lemma** *sum-bounded*:  
**fixes**  $my\text{-seq} :: nat \Rightarrow nat$  **and**  $n :: nat$   
**assumes**  $\bigwedge i. i \geq n \longrightarrow my\text{-seq } i = 0$   
**shows**  $(\sum i. my\text{-seq } i) = (\sum i < n. my\text{-seq } i)$   
 ⟨*proof*⟩

**lemma** *product-seq-eventually-zero*:  
**fixes**  $seq1\ seq2 :: nat \Rightarrow nat$   
**assumes** *eventually-zero*  $seq1$   
**shows** *eventually-zero*  $(\lambda i. seq1\ i * seq2\ i)$   
 ⟨*proof*⟩

**abbreviation** *upper-sum*  
**where**  $upper\text{-sum } seq\ n \equiv \sum i. seq\ (i + n)$   
**syntax**  
 $-from\text{-sum} :: idt \Rightarrow 'a \Rightarrow 'b \Rightarrow 'b \ (\langle (\exists \sum - \geq - / -) \rangle [0,0,10] 10)$   
**syntax-consts**  
 $-from\text{-sum} == upper\text{-sum}$   
**translations**  
 $\sum i \geq n. t == CONST\ upper\text{-sum } (\lambda i. t)\ n$

The following two statements are proved as a sanity check. They are not intended to be used anywhere.

**corollary**  
**fixes**  $seq :: nat \Rightarrow nat$  **and**  $a :: nat$   
**assumes** *seq-def*:  $\bigwedge i. seq\ i = (if\ i = 0\ then\ a\ else\ 0)$   
**shows**  $(\sum i \geq 0. seq\ i) = upper\text{-sum } (\lambda i. seq\ i)\ 0$   
 ⟨*proof*⟩

**corollary**

**fixes**  $seq :: nat \Rightarrow nat$  **and**  $a :: nat$   
**assumes**  $seq-def: \bigwedge i. seq\ i = (if\ i = 0\ then\ a\ else\ 0)$   
**shows**  $(\sum_{i \geq 0}. seq\ i) = a$   
*<proof>*

**lemma** *bounded-sum-from*:

**fixes**  $seq :: nat \Rightarrow nat$  **and**  $n\ s :: nat$   
**assumes**  $\forall i > s. seq\ i = 0$  **and**  $n \leq s$   
**shows**  $(\sum_{i \geq n}. seq\ i) = (\sum_{i=n..s}. seq\ i)$   
*<proof>*

**lemma** *split-suminf*:

**fixes**  $seq :: nat \Rightarrow nat$  **and**  $n :: nat$   
**assumes** *eventually-zero seq*  
**shows**  $(\sum i. seq\ i) = (\sum_{i < n}. seq\ i) + (\sum_{i \geq n}. seq\ i)$   
*<proof>*

**lemma** *dvd-suminf*:

**fixes**  $seq :: nat \Rightarrow nat$  **and**  $b :: nat$   
**assumes** *eventually-zero seq* **and**  $\bigwedge i. b\ dvd\ seq\ i$   
**shows**  $b\ dvd\ (\sum i. seq\ i)$   
*<proof>*

**lemma** *eventually-zero-shifted*:

**assumes** *eventually-zero seq*  
**shows** *eventually-zero*  $(\lambda i. seq\ (i + n))$   
*<proof>*

## 2 Modular arithmetic

This section establishes a number of lemmas about modular arithmetic, including that modular division distributes over an “infinite” sum whose terms are eventually zero.

**lemma** *pmod-int-char*:

**fixes**  $a\ b\ d :: int$   
**shows**  $[a = b] \ (mod\ d) \longleftrightarrow (\exists (n::int). a = b + n*d)$   
*<proof>*

**lemma** *equiv-conj-if*:

**assumes**  $P \Longrightarrow Q$  **and**  $P \Longrightarrow R$  **and**  $Q \Longrightarrow R \Longrightarrow P$   
**shows**  $P \longleftrightarrow Q \wedge R$   
*<proof>*

**lemma** *mod-successor-char*:

**fixes**  $k\ k'\ i\ b :: nat$   
**assumes**  $(b::nat) \geq 2$

**shows**  $[k = k'] \pmod{b^{\wedge}(Suc\ i)} \longleftrightarrow [k\ div\ b^{\wedge}i = k'\ div\ b^{\wedge}i] \pmod{b} \wedge [k = k'] \pmod{b^{\wedge}i}$   
 ⟨proof⟩

**lemma mod-1:**  
**fixes**  $k\ k'\ b :: nat$   
**shows**  $[k = k'] \pmod{b^{\wedge}0}$   
 ⟨proof⟩

**lemma mod-distributes:**  
**fixes**  $seq :: nat \Rightarrow nat$  **and**  $b :: nat$   
**assumes**  $\exists n. \forall i \geq n. seq\ i = 0$   
**shows**  $[(\sum i. seq\ i) = (\sum i. seq\ i\ mod\ b)] \pmod{b}$   
 ⟨proof⟩

**lemma another-mod-cancellation-int:**  
**fixes**  $a\ b\ c\ d\ m :: int$   
**assumes**  $d > 0$  **and**  $[m = a + b] \pmod{c * d}$  **and**  $a\ div\ d = 0$  **and**  $d\ dvd\ b$   
**shows**  $[m\ div\ d = b\ div\ d] \pmod{c}$   
 ⟨proof⟩

**lemma another-mod-cancellation:**  
**fixes**  $a\ b\ c\ d\ m :: nat$   
**assumes**  $d > 0$  **and**  $[m = a + b] \pmod{c * d}$  **and**  $a\ div\ d = 0$  **and**  $d\ dvd\ b$   
**shows**  $[m\ div\ d = b\ div\ d] \pmod{c}$   
 ⟨proof⟩

### 3 Digits as sequence

Rules are introduced for computing the  $i^{\text{th}}$  digit of a base- $b$  representation and the number of digits required to represent a given number. (The latter is essentially an integer version of the base- $b$  logarithm.) It is shown that the sum of the terms  $d_i b^i$  converges to  $m$  if  $d_i$  is the  $i^{\text{th}}$  digit  $m$ . It is shown that the sequence of digits defined is the unique sequence of digits less than  $b$  with this property

Additionally, the `digits_in_base` locale is introduced, which specifies a single symbol  $b$  referring to a natural number greater than one (the base of the digits). Consequently this symbol is omitted from many of the following lemmas and definitions.

**locale digits-in-base =**  
**fixes**  $b :: nat$   
**assumes**  $b\ \text{gte-}2: b \geq 2$   
**begin**

**lemma b-facts [simp]:**  
**shows**  $b > 1$  **and**  $b > 0$  **and**  $b \neq 1$  **and**  $b \neq 0$  **and**  $1\ mod\ b = 1$  **and**  $1\ div\ b = 0$

*<proof>*

Definition based on [1].

**abbreviation** *ith-digit* :: *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *nat* **where**  
*ith-digit* *m* *i*  $\equiv (m \text{ div } b^i) \text{ mod } b$

**lemma** *ith-digit-lt-base*:

**fixes** *m* *i* :: *nat*

**shows**  $0 \leq \text{ith-digit } m \ i$  **and**  $\text{ith-digit } m \ i < b$

*<proof>*

**definition** *num-digits* :: *nat*  $\Rightarrow$  *nat*

**where** *num-digits* *m* = (*LEAST* *i*.  $m < b^i$ )

**lemma** *num-digits-works*:

**shows**  $m < b^{\text{num-digits } m}$

*<proof>*

**lemma** *num-digits-le*:

**assumes**  $m < b^i$

**shows**  $\text{num-digits } m \leq i$

*<proof>*

**lemma** *num-digits-zero*:

**fixes** *m* :: *nat*

**assumes**  $\text{num-digits } m = 0$

**shows**  $m = 0$

*<proof>*

**lemma** *num-digits-gt*:

**assumes**  $m \geq b^i$

**shows**  $\text{num-digits } m > i$

*<proof>*

**lemma** *num-digits-eqI* [*intro*]:

**assumes**  $m \geq b^i$  **and**  $m < b^{i+1}$

**shows**  $\text{num-digits } m = i + 1$

*<proof>*

**lemma** *num-digits-char*:

**assumes**  $m \geq 1$

**shows**  $\text{num-digits } m = i + 1 \iff m \geq b^i \wedge m < b^{i+1}$

*<proof>*

Statement based on [1].

**lemma** *num-digits-recurrence*:

**fixes** *m* :: *nat*

**assumes**  $m \geq 1$

**shows**  $\text{num-digits } m = \text{num-digits } (m \text{ div } b) + 1$

*<proof>*

**lemma** *num-digits-zero-2* [*simp*]: *num-digits 0 = 0*  
*<proof>*

**end**

**locale** *base-10*  
**begin**

As a sanity check, the number of digits in base ten is computed for several natural numbers.

**sublocale** *digits-in-base 10*  
*<proof>*

**corollary**

**shows** *num-digits 0 = 0*  
**and** *num-digits 1 = 1*  
**and** *num-digits 9 = 1*  
**and** *num-digits 10 = 2*  
**and** *num-digits 99 = 2*  
**and** *num-digits 100 = 3*  
*<proof>*

**end**

**context** *digits-in-base*  
**begin**

**lemma** *high-digits-zero-helper*:  
**fixes** *m i :: nat*  
**shows**  $i < \text{num-digits } m \vee \text{ith-digit } m \ i = 0$   
*<proof>*

**lemma** *high-digits-zero*:  
**fixes** *m i :: nat*  
**assumes**  $i \geq \text{num-digits } m$   
**shows**  $\text{ith-digit } m \ i = 0$   
*<proof>*

**lemma** *digit-expansion-bound*:  
**fixes** *i :: nat* **and** *A :: nat  $\Rightarrow$  nat*  
**assumes**  $\bigwedge j. A \ j < b$   
**shows**  $(\sum_{j < i} A \ j * b^{\wedge} j) < b^{\wedge} i$   
*<proof>*

Statement and proof based on [1].

**lemma** *num-digits-suc*:  
**fixes** *n m :: nat*

**assumes**  $Suc\ n = num-digits\ m$   
**shows**  $n = num-digits\ (m\ div\ b)$   
 $\langle proof \rangle$

Proof (and to some extent statement) based on [1].

**lemma** *digit-expansion-bounded-seq*:

**fixes**  $m :: nat$   
**shows**  $m = (\sum_{j < num-digits\ m}. ith-digit\ m\ j * b^{\wedge}j)$   
 $\langle proof \rangle$

A natural number can be obtained from knowing all its base- $b$  digits by the formula  $\sum_j d_j b^j$ .

**theorem** *digit-expansion-seq*:

**fixes**  $m :: nat$   
**shows**  $m = (\sum j. ith-digit\ m\ j * b^{\wedge}j)$   
 $\langle proof \rangle$

**lemma** *lower-terms*:

**fixes**  $a\ c\ i :: nat$   
**assumes**  $c < b^{\wedge}i$  **and**  $a < b$   
**shows**  $ith-digit\ (a * b^{\wedge}i + c)\ i = a$   
 $\langle proof \rangle$

**lemma** *upper-terms*:

**fixes**  $A\ a\ i :: nat$   
**assumes**  $b * b^{\wedge}i\ dvd\ A$  **and**  $a < b$   
**shows**  $ith-digit\ (A + a * b^{\wedge}i)\ i = a$   
 $\langle proof \rangle$

**lemma** *current-term*:

**fixes**  $A\ a\ c\ i :: nat$   
**assumes**  $b * b^{\wedge}i\ dvd\ A$  **and**  $c < b^{\wedge}i$  **and**  $a < b$   
**shows**  $ith-digit\ (A + a * b^{\wedge}i + c)\ i = a$   
 $\langle proof \rangle$

Given that

$$m = \sum_i d_i b^i$$

where the  $d_i$  are all natural numbers less than  $b$ , it follows that  $d_j$  is the  $j^{\text{th}}$  digit of  $m$ .

**theorem** *seq-uniqueness*:

**fixes**  $m\ j :: nat$  **and**  $D :: nat \Rightarrow nat$   
**assumes** *eventually-zero*  $D$  **and**  $m = (\sum i. D\ i * b^{\wedge}i)$  **and**  $\bigwedge i. D\ i < b$   
**shows**  $D\ j = ith-digit\ m\ j$   
 $\langle proof \rangle$

**end**

## 4 Little Endian notation

In this section we begin to define finite digit expansions. Ultimately we want to write digit expansions in “big endian” notation, by which we mean with the highest place digit on the left and the ones digit on the right, since this ordering is standard in informal mathematics. However, it is easier to first define “little endian” expansions with the ones digit on the left since that way the list indexing agrees with the sequence indexing used in the previous section (whenever both are defined).

Notation, definitions, and lemmas in this section typically start with the prefix **LE** (for “little endian”) to distinguish them from the big endian versions in the next section.

**fun** *LEeval-as-base* ( $\langle \_ \_ \rangle$  *LEbase*  $\_$ ) [65, 65] 70)

**where**  $\square$  *LEbase*  $b = 0$   
 $| (d \# \_ \_ \_)$  *LEbase*  $b = d + b * (d\text{-list} \_ \_ \_)$

**corollary** **shows** [2, 4] *LEbase* 5 = (22::nat)  
 $\langle \text{proof} \rangle$

**lemma** *LEbase-one-digit* [simp]: **shows** [ $a::\text{nat}$ ] *LEbase*  $b = a$   
 $\langle \text{proof} \rangle$

**lemma** *LEbase-two-digits* [simp]: **shows** [ $a_0::\text{nat}, a_1$ ] *LEbase*  $b = a_0 + a_1 * b$   
 $\langle \text{proof} \rangle$

**lemma** *LEbase-three-digits* [simp]: **shows** [ $a_0::\text{nat}, a_1, a_2$ ] *LEbase*  $b = a_0 + a_1 * b$   
 $+ a_2 * b^2$   
 $\langle \text{proof} \rangle$

**lemma** *LEbase-closed-form*:  
**shows** ( $A :: \text{nat list}$ ) *LEbase*  $b = (\sum i < \text{length } A . A!i * b^i)$   
 $\langle \text{proof} \rangle$

**lemma** *LEbase-concatenate*:  
**fixes**  $A D :: \text{nat list}$  **and**  $b :: \text{nat}$   
**shows** ( $A @ D$ ) *LEbase*  $b = (A \_ \_ \_)$  +  $b^{\text{length } A} * (D \_ \_ \_)$   
 $\langle \text{proof} \rangle$

**context** *digits-in-base*  
**begin**

**definition** *LEdigits*  $:: \text{nat} \Rightarrow \text{nat list}$  **where**  
*LEdigits*  $m = [\text{ith-digit } m \ i . i \leftarrow [0..<(\text{num-digits } m)]]$

**lemma** *length-is-num-digits*:  
**fixes**  $m :: \text{nat}$   
**shows**  $\text{length } (\text{LEdigits } m) = \text{num-digits } m$

*<proof>*

**lemma** *ith-list-element* [simp]:  
 **assumes**  $(i::nat) < length (LEdigits\ m)$   
 **shows**  $(LEdigits\ m)\ !\ i = ith-digit\ m\ i$   
 *<proof>*

**lemma** *LEbase-infinite-sum*:  
 **fixes**  $m :: nat$   
 **shows**  $(\sum\ i.\ ith-digit\ m\ i * b^i) = (LEdigits\ m)_{LEbase\ b}$   
 *<proof>*

**lemma** *digit-expansion-LElist*:  
 **fixes**  $m :: nat$   
 **shows**  $(LEdigits\ m)_{LEbase\ b} = m$   
 *<proof>*

**lemma** *LElist-uniqueness*:  
 **fixes**  $D :: nat\ list$   
 **assumes**  $\forall\ i < length\ D.\ D!\ i < b$  **and**  $D = [] \vee last\ D \neq 0$   
 **shows**  $LEdigits\ (D_{LEbase\ b}) = D$   
 *<proof>*

**lemma** *LE-digits-zero* [simp]:  $LEdigits\ 0 = []$   
 *<proof>*

**lemma** *LE-units-digit* [simp]:  
 **assumes**  $(m::nat) \in \{1..<b\}$   
 **shows**  $LEdigits\ m = [m]$   
 *<proof>*

**end**

## 5 Big Endian notation

In this section the desired representation of natural numbers, as finite lists of digits with the highest place on the left, is finally realized.

**definition** *BEeval-as-base* ( $\leftarrow_{base}\ \rightarrow$  [65, 65] 70)  
 **where** [simp]:  $D_{base\ b} = (rev\ D)_{LEbase\ b}$

**corollary** **shows**  $[4, 2]_{base\ 5} = (22::nat)$   
 *<proof>*

**lemma** *BEbase-one-digit* [simp]: **shows**  $[a::nat]_{base\ b} = a$   
 *<proof>*

**lemma** *BEbase-two-digits* [simp]: **shows**  $[a_1::nat, a_0]_{base\ b} = a_1*b + a_0$   
 *<proof>*

**lemma** *BEbase-three-digits* [*simp*]: **shows**  $[a_2::nat, a_1, a_0]_{base\ b} = a_2 * b^2 + a_1 * b + a_0$   
 <proof>

**lemma** *BEbase-closed-form*:  
**fixes**  $A :: nat\ list$  **and**  $b :: nat$   
**shows**  $A_{base\ b} = (\sum i < length\ A. A!i * b^{(length\ A - Suc\ i)})$   
 <proof>

**lemma** *BEbase-concatenate*:  
**fixes**  $A\ D :: nat\ list$  **and**  $b :: nat$   
**shows**  $(A @ D)_{base\ b} = (A_{base\ b}) * b^{(length\ D)} + (D_{base\ b})$   
 <proof>

**context** *digits-in-base*  
**begin**

**definition** *digits* ::  $nat \Rightarrow nat\ list$  **where**  
 $digits\ m = rev\ (LEdigits\ m)$

**lemma** *length-is-num-digits-2*:  
**fixes**  $m :: nat$   
**shows**  $length\ (digits\ m) = num-digits\ m$   
 <proof>

**lemma** *LE-BE-equivalence*:  
**fixes**  $m :: nat$   
**shows**  $(digits\ m)_{base\ b} = (LEdigits\ m)_{LEbase\ b}$   
 <proof>

**lemma** *BEbase-infinite-sum*:  
**fixes**  $m :: nat$   
**shows**  $(\sum i. ith-digit\ m\ i * b^i) = (digits\ m)_{base\ b}$   
 <proof>

Every natural number can be represented in base  $b$ , specifically by the digits sequence defined earlier.

**theorem** *digit-expansion-list*:  
**fixes**  $m :: nat$   
**shows**  $(digits\ m)_{base\ b} = m$   
 <proof>

If two natural numbers have the same base- $b$  representation, then they are equal.

**lemma** *digits-cancellation*:  
**fixes**  $k\ m :: nat$   
**assumes**  $digits\ k = digits\ m$   
**shows**  $k = m$

*<proof>*

Suppose we have a finite (possibly empty) sequence  $D_1, \dots, D_n$  of natural numbers such that  $0 \leq D_i < b$  for all  $i$  and such that  $D_1$ , if it exists, is nonzero. Then this sequence is the base- $b$  representation for  $\sum_i D_i b^{n-i}$ .

**theorem** *list-uniqueness*:

**fixes**  $D :: \text{nat list}$

**assumes**  $\forall d \in \text{set } D. d < b$  **and**  $D = [] \vee D!0 \neq 0$

**shows**  $\text{digits } (D_{\text{base } b}) = D$

*<proof>*

We now prove some simplification rules (including a recurrence relation) to make it easier for Isabelle/HOL to compute the base- $b$  representation of a natural number.

The base- $b$  representation of 0 is empty, at least following the conventions of this theory file.

**lemma** *digits-zero [simp]*:

**shows**  $\text{digits } 0 = []$

*<proof>*

If  $0 < m < b$ , then the base- $b$  representation of  $m$  consists of a single digit, namely  $m$  itself.

**lemma** *single-digit-number [simp]*:

**assumes**  $m \in \{0 < .. < b\}$

**shows**  $\text{digits } m = [m]$

*<proof>*

For all  $m \geq b$ , the base- $b$  representation of  $m$  consists of the base- $b$  representation of  $\lfloor m/b \rfloor$  followed by (as the last digit) the remainder of  $m$  when divided by  $b$ .

**lemma** *digits-recurrence [simp]*:

**assumes**  $m \geq b$

**shows**  $\text{digits } m = (\text{digits } (m \text{ div } b)) @ [m \text{ mod } b]$

*<proof>*

**end**

## 6 Exercises

This section contains demonstrations of how to denote certain facts with the notation of the previous sections, and how to quickly prove those facts using the lemmas and theorems above.

The base-5 representation of 22 is  $42_5$ .

**corollary** *digits-in-base*.  $\text{digits } 5 \ 22 = [4, 2]$

*<proof>*

A different proof of the same statement.

**corollary** *digits-in-base.digits* 5 22 = [4, 2]  
*<proof>*

**end**

## References

- [1] B. Porter. Threedivides. <https://isabelle.in.tum.de/dist/library/HOL/HOL-ex/ThreeDivides.html>, 2005. Accessed: 2023-03-06.