

Dedekind Sums

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Abstract

For integers h, k , the Dedekind sum is defined as

$$s(h, k) = \sum_{r=1}^{k-1} \frac{r}{k} \left(\left\{ \frac{hr}{k} \right\} - \frac{1}{2} \right)$$

where $\{x\} = x - [x]$ denotes the fractional part of x .

These sums occur in various contexts in analytic number theory, e.g. in the functional equation of the Dedekind η function or in the study of modular forms.

We give the definition of $s(h, k)$ and prove its basic properties, including the reciprocity law

$$s(h, k) + s(k, h) = \frac{1}{12hk} + \frac{h}{12k} + \frac{k}{12h} - \frac{1}{4}$$

and various congruence results.

Our formalisation follows Chapter 3 of Apostol's *Modular Functions and Dirichlet Series in Number Theory* [1] and contains all facts related to Dedekind sums from it (without the exercises).

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1 Dedekind sums

theory *Dedekind_Sums*

imports

Complex_Main

"HOL-Library.Periodic_Fun"

"HOL-Library.Real_Mod"

"HOL-Number_Theory.Number_Theory"

"Bernoulli.Bernoulli_FPS"

begin

1.1 Preliminaries

lemma *rcong_of_intI*: "[a = b] (mod m) \implies [of_int a = of_int b] (rmod (of_int m))"

by (*metis cong_iff_lin mult.commute of_int_add of_int_mult rcong_altdef*)

lemma *rcong_of_int_iff*: "[of_int a = of_int b] (rmod (of_int m)) \longleftrightarrow [a = b] (mod m)"

proof

assume "[of_int a = of_int b] (rmod (of_int m))"

then obtain *k* where "real_of_int b = of_int a + of_int k * of_int m"

unfolding *rcong_altdef* by *blast*

also have "... = of_int (a + k * m)"

by *simp*

finally have "b = a + k * m"

by *linarith*

thus "[a = b] (mod m)"

by (*simp add: cong_iff_lin*)

qed (*intro rcong_of_intI*)

1.2 Definition and basic properties

definition *dedekind_sum* :: "int \Rightarrow int \Rightarrow real" **where**

"dedekind_sum *h k* $\equiv \sum_{r \in \{1..<k\}} (\text{of_int } r / \text{of_int } k * (\text{frac } (\text{of_int } (h*r) / \text{of_int } k) - 1/2))"$

definition *dedekind_frac* :: "real \Rightarrow real" **where**

"dedekind_frac *x* = (if $x \in \mathbb{Z}$ then 0 else frac $x - 1 / 2$)"

lemma *dedekind_frac_int [simp]*: " $x \in \mathbb{Z} \implies \text{dedekind_frac } x = 0$ "

by (*auto simp: dedekind_frac_def*)

notation *dedekind_frac* (" $\langle _ \rangle$ ")

interpretation *dedekind_frac*: *periodic_fun_simple'* *dedekind_frac*

proof

show " $\langle x + 1 \rangle = \langle x \rangle$ " for *x*

by (*auto simp: dedekind_frac_def frac_1_eq*)

qed

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lemma dedekind_frac_uminus [simp]: " $\langle -x \rangle = -\langle x \rangle$ "
  by (auto simp add: dedekind_frac_def frac_neg)

lemma dedekind_frac_one_minus [simp]: " $\langle 1 - x \rangle = -\langle x \rangle$ "
  by (metis dedekind_frac.minus_1 dedekind_frac_uminus minus_diff_eq)

lemma dedekind_frac_rcong:
  assumes "[ $x = x'$ ] (rmod 1)"
  shows " $\langle x \rangle = \langle x' \rangle$ "
proof -
  from assms obtain k where  $k: "x' = x + \text{of\_int } k"$ 
  by (auto simp: rcong_altdef)
  thus ?thesis
  by (auto simp: dedekind_frac.plus_of_int)
qed

lemma dedekind_frac_mod:
  " $\langle \text{of\_int } (a \bmod k) / \text{of\_int } k \rangle = \langle \text{of\_int } a / \text{of\_int } k \rangle$ "
proof (cases " $k = 0$ ")
  case False
  have "[ $\text{real\_of\_int } (a \bmod k) = \text{real\_of\_int } a$ ] (rmod real_of_int k)"
  by (intro rcong_of_intI cong_mod_leftI cong_refl)
  thus ?thesis using False
  by (intro dedekind_frac_rcong rcong_divide_modulus) auto
qed auto

lemma sum_dedekind_frac_eq_0:
  " $(\sum_{r \in \{1..<k\}} \langle \text{of\_int } r / \text{of\_int } k \rangle) = 0$ "
proof -
  have " $(\sum_{r \in \{1..<k\}} \langle \text{of\_int } (r) / \text{of\_int } k \rangle) =$   

 $(\sum_{r \in \{1..<k\}} \langle \text{of\_int } (k - r) / \text{of\_int } k \rangle)$ "
  by (intro sum.reindex_bij_witness[of _ " $\lambda r. k - r$ " " $\lambda r. k - r$ "])
  auto
  also have " $(\sum_{r \in \{1..<k\}} \langle \text{of\_int } (k - r) / \text{of\_int } k \rangle) =$   

 $(\sum_{r \in \{1..<k\}} \langle 1 + -\text{of\_int } r / \text{of\_int } k \rangle)$ "
  by (intro sum.cong_refl arg_cong[of _ _ dedekind_frac]) (auto simp:
field_simps)
  also have "... =  $-(\sum_{r \in \{1..<k\}} \langle \text{of\_int } r / \text{of\_int } k \rangle)$ "
  by (simp add: sum_negf)
  finally show ?thesis
  by simp
qed

lemma sum_dedekind_aux:
  assumes " $f (0::\text{int}) = 0$ "
  shows " $(\sum_{r \in \{0..<k\}} f r) = (\sum_{r \in \{1..<k\}} f r)$ "
proof (rule sum.mono_neutral_right)
  have " $\{0..<k\} - \{1..<k\} \subseteq \{0\}$ "

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    by auto
  thus " $\forall i \in \{0..<k\} - \{1..<k\}. f\ i = 0$ "
    using assms by blast
qed auto

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lemma sum_dedekind_frac_eq_0':
  " $(\sum_{r \in \{0..<k\}} \langle \text{of\_int } r / \text{of\_int } k \rangle) = 0$ "
proof -
  have " $(\sum_{r \in \{0..<k\}} \langle \text{of\_int } r / \text{of\_int } k \rangle) = (\sum_{r \in \{1..<k\}} \langle \text{of\_int } r / \text{of\_int } k \rangle)$ "
    by (rule sum_dedekind_aux) auto
  with sum_dedekind_frac_eq_0[of k] show ?thesis
    by simp
qed

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lemma sum_dedekind_frac_mult_eq_0:
  assumes "coprime h k"
  shows " $(\sum_{r \in \{1..<k\}} \langle \text{of\_int } (h * r) / \text{of\_int } k \rangle) = 0$ "
proof -
  have " $(\sum_{r \in \{1..<k\}} \langle \text{of\_int } (h * r) / \text{of\_int } k \rangle) = (\sum_{r \in \{1..<k\}} \langle \text{of\_int } ((h * r) \bmod k) / \text{of\_int } k \rangle)$ "
  proof (intro sum.cong)
    fix r assume r: "r  $\in$  {1..<k}"
    have k: "k > 0"
      using r by auto
    have "(h * r) mod k = h * r - k * ((h * r) div k)"
      by (metis minus_mult_div_eq_mod)
    also have "real_of_int ... / of_int k = of_int (h * r) / of_int k - of_int ((h * r) div k)"
      using k by (auto simp: field_simps)
    also have "... =  $\langle \text{of\_int } (h * r) / \text{of\_int } k \rangle$ "
      by (rule dedekind_frac.minus_of_int)
    finally show " $\langle \text{of\_int } (h * r) / \text{of\_int } k \rangle = \langle \text{of\_int } ((h * r) \bmod k) / \text{of\_int } k \rangle$ " ..
  qed auto
  also have "... =  $(\sum_{r \in \{1..<k\}} \langle \text{of\_int } r / \text{of\_int } k \rangle)$ "
    by (rule sum.reindex_bij_betw[OF bij_betw_int_remainders_mult[OF assms]])
  also have "... = 0"
    by (rule sum_dedekind_frac_eq_0)
  finally show ?thesis .
qed

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lemma sum_dedekind_frac_mult_eq_0':
  assumes "coprime h k"
  shows " $(\sum_{r \in \{0..<k\}} \langle \text{of\_int } (h * r) / \text{of\_int } k \rangle) = 0$ "
proof -
  have " $(\sum_{r \in \{0..<k\}} \langle \text{of\_int } (h * r) / \text{of\_int } k \rangle) = (\sum_{r \in \{1..<k\}} \langle \text{of\_int } (h * r) / \text{of\_int } k \rangle)$ "
    by (intro sum_dedekind_aux) auto

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also have "... = 0"
  by (intro sum_dedekind_frac_mult_eq_0 assms)
finally show ?thesis .
qed

lemma dedekind_sum_altdef:
  assumes "coprime h k"
  shows "dedekind_sum h k = ( $\sum_{r \in \{1..<k\}}$ .  $\langle \text{of\_int } r / \text{of\_int } k \rangle * \langle \text{of\_int } (h*r) / \text{of\_int } k \rangle$ )"
proof -
  have " $(\sum_{r \in \{1..<k\}}$ .  $\langle \text{of\_int } r / \text{of\_int } k \rangle * \langle \text{of\_int } (h*r) / \text{of\_int } k \rangle) =$ 
    ( $\sum_{r \in \{1..<k\}}$ .  $(\text{of\_int } r / \text{of\_int } k - 1 / 2) * \langle \text{of\_int } (h*r) / \text{of\_int } k \rangle$ )"
  proof (intro sum.cong ballI)
    fix r assume r: "r  $\in$  {1..<k}"
    hence "1  $\leq$  r" "r < k"
    by auto
    hence " $\neg k \text{ dvd } r$ "
    using zdvd_not_zless by auto
    hence "real_of_int r / real_of_int k  $\notin$   $\mathbb{Z}$ "
    using r by (subst of_int_div_of_int_in_Ints_iff) auto
    moreover have "frac (real_of_int r / real_of_int k) = real_of_int r / real_of_int k"
    using r by (subst frac_eq) auto
    ultimately show " $\langle \text{real\_of\_int } r / \text{real\_of\_int } k \rangle * \langle \text{real\_of\_int } (h * r) / \text{real\_of\_int } k \rangle =$ 
      ( $\text{real\_of\_int } r / \text{real\_of\_int } k - 1 / 2) * \langle \text{real\_of\_int } (h * r) / \text{real\_of\_int } k \rangle$ "
    by (subst dedekind_frac_def) auto
  qed auto
  also have "... = ( $\sum_{r \in \{1..<k\}}$ .  $\text{of\_int } r / \text{of\_int } k * \langle \text{of\_int } (h*r) / \text{of\_int } k \rangle$ ) -
    1 / 2 * ( $\sum_{r \in \{1..<k\}}$ .  $\langle \text{of\_int } (h*r) / \text{of\_int } k \rangle$ )"
  unfolding sum_distrib_left sum_subtractf [symmetric] by (simp add: ring_distrib)
  also have " $(\sum_{r \in \{1..<k\}}$ .  $\langle \text{of\_int } (h*r) / \text{of\_int } k \rangle) = 0$ "
  by (intro sum_dedekind_frac_mult_eq_0 assms)
  also have " $(\sum_{r \in \{1..<k\}}$ .  $\text{of\_int } r / \text{of\_int } k * \langle \text{of\_int } (h*r) / \text{of\_int } k \rangle) = \text{dedekind\_sum } h \text{ } k$ "
  unfolding dedekind_sum_def
proof (intro sum.cong)
  fix r assume r: "r  $\in$  {1..<k}"
  have " $\neg k \text{ dvd } h * r$ "
  proof
    assume "k dvd h * r"
    with assms have "k dvd r"
    using coprime_commute coprime_dvd_mult_right_iff by blast
    hence "k  $\leq$  r"
  qed

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    using r by (intro zdvd_imp_le) auto
  thus False
    using r by simp
qed
hence "real_of_int (h * r) / real_of_int k  $\notin$   $\mathbb{Z}$ "
  using r by (subst of_int_div_of_int_in_Ints_iff) auto
thus "real_of_int r / real_of_int k *  $\langle$ real_of_int (h * r) / real_of_int
k $\rangle$  =
      real_of_int r / real_of_int k * (frac (real_of_int (h * r) /
real_of_int k) - 1 / 2)"
  by (auto simp: dedekind_frac_def)
qed auto
finally show ?thesis
  by simp
qed

theorem dedekind_sum_cong:
  assumes "[h' = h] (mod k)"
  assumes "coprime h' k  $\vee$  coprime h k"
  shows "dedekind_sum h' k = dedekind_sum h k"
proof -
  have coprime: "coprime h' k" "coprime h k"
    using coprime_cong_cong_left[OF assms(1)] assms(2) by auto
  show ?thesis
    unfolding coprime[THEN dedekind_sum_altdef]
  proof (intro sum.cong)
    fix r assume r: "r  $\in$  {1.. $k$ }"
    have "[real_of_int (h' * r) = real_of_int (h * r)] (rmod real_of_int
k)"
      by (intro rcong_of_intI cong_mult cong_refl assms)
    hence " $\langle$ real_of_int (h' * r) / real_of_int k $\rangle$  =  $\langle$ real_of_int (h *
r) / real_of_int k $\rangle$ "
      using r by (intro dedekind_frac_rcong rcong_divide_modulus) auto
    thus " $\langle$ real_of_int r / real_of_int k $\rangle$  *  $\langle$ real_of_int (h' * r) / real_of_int
k $\rangle$  =
           $\langle$ real_of_int r / real_of_int k $\rangle$  *  $\langle$ real_of_int (h * r) / real_of_int
k $\rangle$ "
      by simp
    qed auto
  qed
qed

theorem dedekind_sum_negate:
  assumes "coprime h k"
  shows "dedekind_sum (-h) k = -dedekind_sum h k"
proof -
  have *: "coprime (-h) k"
    using assms by auto
  show ?thesis

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    unfolding dedekind_sum_altdef[OF assms] dedekind_sum_altdef[OF *]
    by (simp_all add: sum_negf)
qed

theorem dedekind_sum_negate_cong:
  assumes "[h' = -h] (mod k)" "coprime h' k  $\vee$  coprime h k"
  shows "dedekind_sum h' k = -dedekind_sum h k"
proof -
  have coprime: "coprime h' k" "coprime h k"
    using coprime_cong_cong_left[OF assms(1)] assms(2) by auto
  from coprime show ?thesis
    using assms(1) dedekind_sum_cong dedekind_sum_negate by metis
qed

theorem dedekind_sum_inverse:
  assumes "[h * h' = 1] (mod k)"
  shows "dedekind_sum h k = dedekind_sum h' k"
proof -
  have 1: "coprime h k"
    using assms coprime_iff_invertible_int by blast
  have 2: "coprime h' k"
    using assms coprime_iff_invertible_int by (subst (asm) mult.commute)
auto
  have "dedekind_sum h' k =
    ( $\sum r = 1..<k. \langle \text{real\_of\_int } (1 * r) / \text{real\_of\_int } k \rangle * \langle \text{real\_of\_int } (h' * r) / \text{real\_of\_int } k \rangle$ )"
    unfolding dedekind_sum_altdef[OF 2] by simp
  also have "... = ( $\sum r = 1..<k. \langle \text{real\_of\_int } (h * ((h' * r) \bmod k)) / \text{real\_of\_int } k \rangle * \langle \text{real\_of\_int } ((h' * r) \bmod k) / \text{real\_of\_int } k \rangle$ )"
    proof (rule sum.cong, goal_cases)
      case (2 r)
      have "[real_of_int (1 * r) = real_of_int (h * h' * r)] (rmod real_of_int k)"
        using assms by (intro rcong_of_intI cong_mult cong_refl assms) (auto simp: cong_sym)
      also have "[real_of_int (h * h' * r) = real_of_int (h * ((h' * r) mod k))] (rmod of_int k)"
        by (subst mult.assoc, intro rcong_of_intI cong_mult cong_refl cong_mod_rightI)
      finally have *: " $\langle \text{real\_of\_int } (1 * r) / \text{real\_of\_int } k \rangle = \langle \text{real\_of\_int } (h * ((h' * r) \bmod k)) / \text{real\_of\_int } k \rangle$ "
        using 2 by (intro ext dedekind_frac_rcong rcong_divide_modulus)
    auto
  have "[real_of_int (h' * r) = real_of_int (h' * r mod k)] (rmod real_of_int k)"
    by (intro rcong_of_intI cong_refl cong_mod_rightI)
  hence " $\langle \text{real\_of\_int } (h' * r) / \text{real\_of\_int } k \rangle = \langle \text{real\_of\_int } ((h' * r) \bmod k) / \text{real\_of\_int } k \rangle$ "

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* r) mod k) / real_of_int k)"
  using 2 by (intro dedekind_frac_rcong rcong_divide_modulus) auto
  thus ?case using *
  by (simp add: mult_ac)
qed auto
also have "... = ( $\sum r = 1..<k. \langle \text{real\_of\_int } (h * r) / \text{real\_of\_int } k \rangle$ )
*  $\langle \text{real\_of\_int } r / \text{real\_of\_int } k \rangle$ )"
  by (rule sum.reindex_bij_betw [OF bij_betw_int_remainders_mult]) fact
also have "... = dedekind_sum h k"
  using 1 by (simp add: dedekind_sum_altdef mult_ac)
finally show ?thesis ..
qed

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theorem dedekind_sum_inverse':
  assumes "[h * h' = -1] (mod k)"
  shows "dedekind_sum h k = -dedekind_sum h' k"
proof -
  have "[-(h * h') = - (-1)] (mod k)"
    using assms unfolding cong_minus_minus_iff .
  hence 1: "[h * -h' = 1] (mod k)"
    by simp
  hence "[h' * (-h) = 1] (mod k)"
    by (simp add: mult_ac)
  hence 2: "coprime h' k"
    using assms coprime_iff_invertible_int by blast
  from 1 have "dedekind_sum h k = dedekind_sum (-h') k"
    by (intro dedekind_sum_inverse)
  thus ?thesis
    using 2 by (simp add: dedekind_sum_negate)
qed

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theorem dedekind_sum_eq_zero:
  assumes "[h2 + 1 = 0] (mod k)"
  shows "dedekind_sum h k = 0"
proof -
  have "[h * h = h2 + 1 - 1] (mod k)"
    by (simp add: algebra_simps power2_eq_square)
  also have "[h2 + 1 - 1 = 0 - 1] (mod k)"
    by (intro cong_diff assms cong_refl)
  finally have "[h * h = -1] (mod k)"
    by simp
  hence "dedekind_sum h k = -dedekind_sum h k"
    by (rule dedekind_sum_inverse')
  thus ?thesis
    by simp
qed

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1.3 The Reciprocity Law

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theorem sum_of_powers':
  "(\sum k<n::nat. (real k) ^ m) = (bernpoly (Suc m) n - bernpoly (Suc m)
  0) / (m + 1)"
proof (cases "n = 0")
  case True
  thus ?thesis
    by auto
next
  case False
  hence "{..<n} = {...n-1}"
    by auto
  thus ?thesis
    using sum_of_powers[of m "n - 1"] False by (simp add: of_nat_diff)
qed

theorem sum_of_powers'_int:
  assumes "n ≥ 0"
  shows "(\sum k=0..<n::int. real_of_int k ^ m) = (bernpoly (Suc m) n
  - bernpoly (Suc m) 0) / (m + 1)"
proof -
  have "(\sum k=0..<n::int. real_of_int k ^ m) = (\sum k<nat n::nat. real k
  ^ m)"
    by (intro sum.reindex_bij_witness[of _ int nat]) auto
  also have "(\sum k<nat n::nat. real k ^ m) = (bernpoly (Suc m) n - bernpoly
  (Suc m) 0) / (m + 1)"
    using assms by (subst sum_of_powers') (auto)
  finally show ?thesis by simp
qed

theorem sum_of_powers'_int_from_1:
  assumes "n ≥ 0" "m > 0"
  shows "(\sum k=1..<n::int. real_of_int k ^ m) = (bernpoly (Suc m) n
  - bernpoly (Suc m) 0) / (m + 1)"
proof -
  have "(\sum k=1..<n::int. real_of_int k ^ m) = (\sum k=0..<n::int. real_of_int
  k ^ m)"
    using assms by (intro sum.mono_neutral_left) auto
  also have "... = (bernpoly (Suc m) n - bernpoly (Suc m) 0) / (m + 1)"
    using assms by (subst sum_of_powers'_int) auto
  finally show ?thesis by simp
qed

theorem dedekind_sum_reciprocity:
  assumes "h > 0" and "k > 0" and "coprime h k"
  shows "12 * h * k * dedekind_sum h k + 12 * k * h * dedekind_sum k h
  =
  h^2 + k^2 - 3 * h * k + 1"

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proof -
  have upto_3_eq [simp]: "{..(3::nat)} = {0, 1, 2, 3}"
    by auto
  have [simp]: "(3 choose 2) = 3"
    by (simp add: eval_nat_numeral)

  define S where "S = ( $\sum_{r \in \{1..<k\}} \langle \text{of\_int } (h*r) / \text{of\_int } k \rangle ^ 2$ )"
  define T where "T = ( $\sum_{r = 1..<k} \lfloor \text{real\_of\_int } (h * r) / k \rfloor * (\lfloor \text{real\_of\_int } (h * r) / k \rfloor + 1)$ )"

  have "S = ( $\sum_{r \in \{1..<k\}} \langle \text{of\_int } ((h * r) \bmod k) / \text{of\_int } k \rangle ^ 2$ )"
    unfolding S_def
    by (intro sum.cong arg_cong[of _ _ "\lambda x. x ^ 2"] dedekind_frac_mod
    [symmetric] refl)
  also have "... = ( $\sum_{r \in \{1..<k\}} \langle \text{of\_int } r / \text{of\_int } k \rangle ^ 2$ )"
    by (rule sum.reindex_bij_betw[OF bij_betw_int_remainders_mult]) fact
  also have "... = ( $\sum_{r \in \{1..<k\}} (\text{of\_int } r / \text{of\_int } k - 1 / 2) ^ 2$ )"
  proof (rule sum.cong, goal_cases)
    case (2 r)
    have "~k dvd r"
      using 2 by (auto dest!: zdvd_imp_le)
    hence "of_int r / real_of_int k  $\notin \mathbb{Z}$ "
      using 2 by (subst of_int_div_of_int_in_Ints_iff) auto
    moreover have "frac (of_int r / real_of_int k) = of_int r / of_int
    k"
      using 2 by (subst frac_eq) auto
    ultimately show ?case
      by (simp add: dedekind_frac_def)
  qed auto
  also have "... = 1 / k2 * ( $\sum_{r=1..<k} \text{of\_int } r ^ 2$ ) -
    1 / k * ( $\sum_{r=1..<k} \text{of\_int } r$ ) +
    ( $\sum_{r=1..<k} 1 / 4$ )"
    unfolding sum_distrib_left sum_subtractf [symmetric] sum.distrib [symmetric]
    dedekind_sum_def
    by (intro sum.cong) (auto simp: field_simps power2_eq_square)
  finally have "S = ..." .

  note this [symmetric]
  also have "S = ( $\sum_{r = 1..<k} (\text{frac } (\text{real\_of\_int } (h * r) / \text{real\_of\_int } k) - 1 / 2)^2$ )"
    unfolding S_def
  proof (rule sum.cong, goal_cases)
    case (2 r)
    have "~k dvd r"
      using 2 by (auto dest!: zdvd_imp_le)
    hence "~k dvd h * r"
      using assms(3) coprime_commute coprime_dvd_mult_right_iff by blast
    hence "of_int (h * r) / real_of_int k  $\notin \mathbb{Z}$ "
      using 2 by (subst of_int_div_of_int_in_Ints_iff) auto
  end

```

```

thus ?case
  by (simp add: dedekind_frac_def)
qed auto
also have "... =
  2 * h * dedekind_sum h k +
  (∑ r=1..<k. real_of_int ([of_int(h*r)/k] * ([h*r/k] + 1))) -
  h2/k2 * (∑ r=1..<k. real_of_int r ^ 2) +
  (∑ r=1..<k. 1/4)"
  unfolding sum_distrib_left sum_subtractf [symmetric] sum.distrib [symmetric]
dedekind_sum_def
  by (intro sum.cong) (auto simp: field_simps power2_eq_square frac_def)
  finally have "2 * h * dedekind_sum h k =
    -(∑ r=1..<k. real_of_int ([of_int(h*r)/k] * ([h*r/k]
+ 1))) +
    (h2 + 1)/k2 * (∑ r=1..<k. real_of_int r ^ 2) - 1/k *
(∑ r=1..<k. of_int r)"
  by (simp add: add_divide_distrib ring_distrib)
  hence "6 * k * (2 * h * dedekind_sum h k) = 6 * k * (
    -(∑ r=1..<k. real_of_int ([of_int(h*r)/k] * ([h*r/k] + 1)))
+
    (h2 + 1)/k2 * (∑ r=1..<k. real_of_int r ^ 2) - 1/k * (∑ r=1..<k.
of_int r))"
  by (simp only: )
  also have "... = -6 * k * (∑ r=1..<k. real_of_int ([of_int(h*r)/k] *
([h*r/k] + 1))) +
    6 * (h2 + 1)/k * (∑ r=1..<k. real_of_int r ^ 2) -
    6 * (∑ r=1..<k. real_of_int r)"
  using <k > 0> by (simp add: field_simps power2_eq_square)
  also have "(∑ v=1..<k. real_of_int v) = k ^ 2 / 2 - k / 2"
  using sum_of_powers'_int_from_1[of k 1] assms
  by (simp add: bernpoly_def algebra_simps power2_eq_square)
  also have "6 * ... = 3 * k2 - 3 * k"
  using assms by (simp add: field_simps)
  also have "(∑ v=1..<k. real_of_int v ^ 2) = k ^ 3 / 3 - k ^ 2 / 2 +
k / 6"
  using assms by (subst sum_of_powers'_int_from_1) (auto simp: bernpoly_def)
  also have "real_of_int (6 * (h2 + 1)) / real_of_int k * ... =
    2 * h2 * k2 + 2 * k2 - 3 * h2 * k - 3 * k + h2 + 1"
  using assms by (simp add: field_simps power3_eq_cube power2_eq_square)
  finally have sum_eq1: "12 * h * k * dedekind_sum h k =
    -6 * k * T + 2 * h2 * k2 + 2 * k2 - 3 * h2 * k - 3 * k + h2 + 1 -
3 * k2 + 3 * k"
  by (simp add: mult_ac T_def)

define N where "N = (λv. card {r∈{1..<k}. [real_of_int (h*r)/k] = v
- 1})"
have N_eq_aux: "{r∈{1..<k}. [real_of_int (h*r)/k] = v - 1} =
  {1..<k} ∩ {[k * (v - 1) / h] <.. [k * v / h]}" for v
proof -

```

```

    have "[real_of_int (h*r)/k] = v - 1  $\longleftrightarrow$  r  $\in$  {[k * (v - 1) / h]<.. $\lfloor$ k
* v / h $\rfloor$ }"
    if r: "r  $\in$  {1.. $\lfloor$ k}" for r
  proof -
    have neq: "real_of_int r  $\neq$  k * v / h" for v
    proof
      assume "real_of_int r = k * v / h"
      hence "real_of_int (r * h) = real_of_int (k * v)"
        using assms unfolding of_int_mult by (simp add: field_simps)
      hence *: "r * h = k * v"
        by linarith
      have "k dvd r * h"
        by (subst *) auto
      with <coprime h k> have "k dvd r"
        by (simp add: coprime_commute coprime_dvd_mult_left_iff)
      with r show False
        by (auto dest!: zdvd_imp_le)
    qed
  qed

  have "[real_of_int (h*r)/k] = v - 1  $\longleftrightarrow$  h * r / k  $\in$  {v-1.. $\lfloor$ v}"
    unfolding atLeastLessThan_iff by linarith
  also have "...  $\longleftrightarrow$  r  $\in$  {k * (v - 1) / h.. $\lfloor$ k * v / h}"
    using assms by (auto simp: field_simps)
  also have "...  $\longleftrightarrow$  r  $\in$  {k * (v - 1) / h<.. $\lfloor$ k * v / h}"
    using neq[of "v - 1"] neq[of v] by auto
  also have "...  $\longleftrightarrow$  r  $\in$  {[k * (v - 1) / h]<.. $\lfloor$ k * v / h $\rfloor$ }"
    unfolding greaterThanAtMost_iff by safe linarith+
  finally show "[real_of_int (h*r)/k] = v - 1  $\longleftrightarrow$  r  $\in$  {[k * (v -
1) / h]<.. $\lfloor$ k * v / h $\rfloor$ }" .
  qed
  thus "{r $\in$ {1.. $\lfloor$ k}. [real_of_int (h*r)/k] = v - 1} =
    {1.. $\lfloor$ k}  $\cap$  {[k * (v - 1) / h]<.. $\lfloor$ k * v / h $\rfloor$ }"
    by blast
  qed

define N' where "N' = ( $\lambda$ v. if v = h then k - 1 else  $\lfloor$ k * v / h $\rfloor$ )"

have N_eq: "int (N v) = N' v - N' (v - 1)" if v: "v  $\in$  {1..h}" for v
proof (cases "v = h")
  case False
    have le: "[real_of_int k * (real_of_int v - 1) / real_of_int h]
       $\leq$  [real_of_int k * real_of_int v / real_of_int h]"
      using assms by (intro floor_mono divide_right_mono mult_left_mono)
  auto
  have "N v = card ({1.. $\lfloor$ k}  $\cap$  {[k * (v - 1) / h]<.. $\lfloor$ k * v / h $\rfloor$ })"
    unfolding N_def N_eq_aux ..
  also have "{[k * (v - 1) / h]<.. $\lfloor$ k * v / h $\rfloor$ }  $\subseteq$  {1.. $\lfloor$ k}"
  proof -
    have "[k * (v - 1) / h]  $\geq$  0"

```

```

        using v <k > 0> <h > 0> by auto
        moreover have "[k * v / h] < k"
        using v False <k > 0> <h > 0> by (subst floor_less_iff) (auto
simp: field_simps)
        ultimately have "{[k * (v - 1) / h]<..[k * v / h]} ⊆ {0<..k-1}"
        unfolding Ioc_subset_iff by auto
        also have "{0<..k-1} = {1..<k}"
        by force
        finally show ?thesis .
    qed
    hence "{1..<k} ∩ {[k * (v - 1) / h]<..[k * v / h]} = {[k * (v - 1)
/ h]<..[k * v / h]}"
    by blast
    finally show ?thesis
    using le v False by (simp add: N'_def)
next
    case [simp]: True
    have le: "[real_of_int k * (real_of_int h - 1) / real_of_int h] ≤
k - 1"
    using assms by (subst floor_le_iff) (auto simp: field_simps)
    have "N h = card ({1..<k} ∩ {[k * (h - 1) / h]<..[k * h / h]})"
    unfolding N_def N_eq_aux ..
    also have "[k * h / h] = k"
    using assms by simp
    also have "{1..<k} ∩ {[real_of_int (k * (h - 1)) / real_of_int h]<..k}
=
        {[real_of_int (k * (h - 1)) / real_of_int h]+1..<k}"
    proof -
        have nonneg: "[real_of_int k * (real_of_int h - 1) / real_of_int
h] ≥ 0"
        unfolding of_int_mult using assms by auto
        have "{1..<k} ∩ {[real_of_int (k * (h - 1)) / real_of_int h]+1..<k+1}
=
            {max 1 ([real_of_int k * (real_of_int h - 1) / real_of_int
h]+1)..<k}"
        by simp
        also have "max 1 ([real_of_int k * (real_of_int h - 1) / real_of_int
h]+1) =
            [real_of_int k * (real_of_int h - 1) / real_of_int h]
+ 1"
        using nonneg by auto
        also have "{[real_of_int (k * (h - 1)) / real_of_int h]+1..<k+1}
=
            {[real_of_int (k * (h - 1)) / real_of_int h]<..k}"
        by force
        finally show ?thesis by simp
    qed
    finally show ?thesis
    using le by (simp add: N'_def)

```

```

qed

have "T = ( $\sum (v,r) \in (\text{SIGMA } v:\{1..h\}. \{r \in \{1..<k\}. \lfloor \text{of\_int } (h*r)/k \rfloor = v - 1\}) . (v - 1) * v)$ "
  unfolding T_def
proof (intro sum.reindex_bij_witness[of _ snd "\lambda r. (\lfloor \text{of\_int } (h*r)/k \rfloor + 1, r)"], goal_cases)
  case (2 r)
  have "\lfloor \text{real\_of\_int } h * \text{real\_of\_int } r / \text{real\_of\_int } k \rfloor < h"
    using 2 assms by (subst floor_less_iff) (auto simp: field_simps)
  thus ?case using 2 assms by auto
qed (use assms in auto)
also have "... = ( $\sum v=1..h. \sum r/r \in \{1..<k\} \wedge \lfloor \text{real\_of\_int}(h*r)/k \rfloor = v - 1. (v - 1) * v$ )"
proof (intro sum.Sigma [symmetric] ballI)
  show "finite {r \in \{1..<k\}. \lfloor \text{real\_of\_int } (h * r) / \text{real\_of\_int } k \rfloor = v - 1}" for v
    by (rule finite_subset[of _ "{1..<k}"]) auto
qed auto
also have "... = ( $\sum v=1..h. (v - 1) * v * \text{int } (N v)$ )"
  by (simp add: N_def mult_ac)
also have "... = ( $\sum v=1..h. (v - 1) * v * (N' v - N' (v - 1))$ )"
  by (intro sum.cong) (auto simp: N_eq)
also have "... = ( $\sum v=1..h. (v - 1) * v * N' v - \sum v=1..h. (v - 1) * v * N' (v - 1)$ )"
  by (simp add: ring_distrib sum_subtractf)
also have "( $\sum v=1..h. (v - 1) * v * N' (v - 1)$ ) = ( $\sum v=0..<h. (v + 1) * v * N' v$ )"
  by (rule sum.reindex_bij_witness[of _ "\lambda v. v+1" "\lambda v. v-1"]) auto
also have "... = ( $\sum v=1..<h. (v + 1) * v * N' v$ )"
  by (intro sum.mono_neutral_right) auto
also have "{1..h} = insert h {1..<h}"
  using assms by auto
also have "( $\sum v \in \dots. (v - 1) * v * N' v$ ) = ( $\sum v = 1..<h. (v - 1) * v * N' v$ ) + (h - 1) * h * N' h"
  by (subst sum.insert) auto
also have "( $\sum v=1..<h. (v-1) * v * N' v$ ) + (h-1) * h * N' h - ( $\sum v=1..<h. (v+1) * v * N' v$ ) = (h-1) * h * N' h - 2 * ( $\sum v=1..<h. v * N' v$ )"
  by (simp add: ring_distrib sum_subtractf sum.distrib sum_distrib_left sum_negf mult_ac)
also have "( $\sum v = 1..<h. v * N' v$ ) = ( $\sum v = 1..<h. v * \lfloor \text{of\_int } (k * v)/h \rfloor$ )"
  by (intro sum.cong) (auto simp: N'_def)
also have "N' h = k - 1"
  by (simp add: N'_def)
finally have *: "- 2 * ( $\sum v = 1..<h. v * \lfloor \text{of\_int } (k * v)/h \rfloor$ ) = -(h - 1) * h * (k - 1) + T"
  by linarith

```

```

have "12 * k * h * dedekind_sum k h =
  6 * k * (-2 * (∑ v=1..<h. v * ⌊real_of_int (k * v)/h⌋)) +
  12*k^2/h * (∑ v=1..<h. real_of_int v ^ 2) -
  6 * k * (∑ v=1..<h. real_of_int v)"
  by (simp add: dedekind_sum_def algebra_simps power2_eq_square sum.distrib
sum_subtractf
      sum_distrib_left sum_distrib_right sum_negf frac_def
sum_divide_distrib)
  also note *
  also have "real_of_int (6 * k * (- (h - 1) * h * (k - 1) + T)) =
  6 * k * T - 6 * h^2 * k^2 + 6 * h * k^2 + 6 * h^2 * k - 6 * h
* k"
  by (simp add: algebra_simps power2_eq_square)
  also have "(∑ v=1..<h. real_of_int v) = h ^ 2 / 2 - h / 2"
  using sum_of_powers'_int_from_1[of h 1] assms
  by (simp add: bernpoly_def algebra_simps power2_eq_square)
  also have "real_of_int (6 * k) * ... = 3 * h^2 * k - 3 * h * k"
  using assms by (simp add: field_simps)
  also have "(∑ v=1..<h. real_of_int v ^ 2) = h ^ 3 / 3 - h ^ 2 / 2 +
h / 6"
  using assms by (subst sum_of_powers'_int_from_1) (auto simp: bernpoly_def)
  also have "real_of_int (12 * k^2) / real_of_int h * ... =
  4 * h^2 * k^2 - 6 * h * k^2 + 2 * k^2"
  using assms by (simp add: field_simps power3_eq_cube power2_eq_square)
  finally have sum_eq2: "12 * k * h * dedekind_sum k h =
  6 * k * T - 2 * h^2 * k^2 + 3 * h^2 * k - 3 * h
* k + 2 * k^2"
  by simp

have "real_of_int (12 * h * k) * dedekind_sum h k + real_of_int (12
* k * h) * dedekind_sum k h =
  real_of_int (h^2 - 3 * h * k + k^2 + 1)"
  unfolding sum_eq1 sum_eq2 by simp
  thus ?thesis
  by simp
qed

theorem dedekind_sum_reciprocity':
  assumes "h > 0" and "k > 0" and "coprime h k"
  shows "dedekind_sum h k = -dedekind_sum k h + h / k / 12 + k / h / 12
- 1 / 4 + 1 / (12 * h * k)"
  using dedekind_sum_reciprocity[OF assms] assms
  by (auto simp: field_simps power2_eq_square)

```

1.4 Congruence Properties

definition `dedekind_sum'` :: "int ⇒ int ⇒ int" where

```

"dedekind_sum' h k = ⌊6 * real_of_int k * dedekind_sum h k⌋"

lemma dedekind_sum'_cong:
  "[h = h'] (mod k)  $\implies$  coprime h k  $\vee$  coprime h' k  $\implies$  dedekind_sum'
  h k = dedekind_sum' h' k"
  unfolding dedekind_sum'_def by (subst dedekind_sum_cong[of h h' k])
  auto

lemma
  assumes "k > 0"
  shows of_int_dedekind_sum':
    "real_of_int (dedekind_sum' h k) = 6 * real_of_int k * dedekind_sum
  h k"
  and dedekind_sum'_altdef:
    "dedekind_sum' h k = h * (k - 1) * (2 * k - 1) -
    6 * ( $\sum_{r=1..<k. r * \lfloor \text{of\_int } (h * r) / k \rfloor$ ) - 3 * (k *
  (k - 1) div 2)"
  and dedekind_sum'_cong_3: "[dedekind_sum' h k = h * (k - 1) * (2
  * k - 1)] (mod 3)"
  proof -
    have [simp]: "{..3} = {0,1,2,3::nat}"
      by auto
    have [simp]: "(3 choose 2) = 3"
      by (auto simp: eval_nat_numeral)
    have "6 * k * dedekind_sum h k = 6 * h / k * ( $\sum_{r=1..<k. \text{of\_int } r ^
  2$ ) -
    6 * ( $\sum_{r=1..<k. \text{of\_int } r * \lfloor h*r/k \rfloor$ ) - 3 * ( $\sum_{r=1..<k. \text{of\_int }
  r ^ 1$ )"
      by (simp add: dedekind_sum_def frac_def algebra_simps sum_distrib_left
  sum_distrib_right
  sum_subtractf sum.distrib sum_divide_distrib power2_eq_square)
    also have "6 * h / k * ( $\sum_{r=1..<k. \text{real\_of\_int } r ^ 2$ ) = 2 * h * k2 +
  h - 3 * h * k"
      using assms by (subst sum_of_powers'_int_from_1)
      (auto simp: bernpoly_def field_simps power2_eq_square
  power3_eq_cube)
    also have "3 * ( $\sum_{r=1..<k. \text{real\_of\_int } r ^ 1$ ) = 3 * (k * (k - 1) / 2)"
      using assms by (subst sum_of_powers'_int_from_1)
      (auto simp: bernpoly_def field_simps power2_eq_square)
    also have "even (k * (k - 1))"
      by auto
    hence "(k * (k - 1) / 2) = real_of_int ((k * (k - 1)) div 2)"
      by fastforce
    also have "2 * h * k2 + h - 3 * h * k = h * (k - 1) * (2 * k - 1)"
      by (simp add: algebra_simps power2_eq_square)
    finally have eq: "6 * real_of_int k * dedekind_sum h k =
    real_of_int (h * (k - 1) * (2 * k - 1) -
    6 * ( $\sum_{r=1..<k. r * \lfloor \text{of\_int } (h * r) / k \rfloor$ )) - 3

```

```

* (k * (k - 1) div 2))"
  unfolding of_int_diff of_int_add by simp

  have "6 * real_of_int k * dedekind_sum h k ∈ ℤ"
    unfolding eq by (rule Ints_of_int)
  thus "real_of_int (dedekind_sum' h k) = 6 * of_int k * dedekind_sum
h k"
    unfolding dedekind_sum'_def by (auto elim!: Ints_cases)
  show eq': "dedekind_sum' h k = h * (k - 1) * (2 * k - 1) -
6 * (∑ r = 1..<k. r * [of_int (h * r) / k]) - 3 * (k *
(k - 1) div 2)"
    unfolding dedekind_sum'_def eq by (simp only: floor_of_int)
  have "[dedekind_sum' h k = h * (k - 1) * (2 * k - 1) - 0 - 0] (mod 3)"
    unfolding eq' by (intro cong_diff cong_refl) (auto simp: Cong.cong_def)
  thus "[dedekind_sum' h k = h * (k - 1) * (2 * k - 1)] (mod 3)"
    by simp
qed

lemma three_dvd_dedekind_sum'_iff_aux:
  fixes h k :: int
  defines "∅ ≡ gcd 3 k"
  assumes "k > 0" "coprime h k"
  shows "3 dvd (2 * dedekind_sum' h k) ⟷ ¬3 dvd k"
proof -
  have "[2 * dedekind_sum' h k = 2 * (h * (k - 1) * (2 * k - 1))] (mod
3)"
    by (intro cong_mult dedekind_sum'_cong_3 cong_refl assms)
  also have "2 * (h * (k - 1) * (2 * k - 1)) = h * (k - 1) * (4 * k +
(-2))"
    by (simp add: algebra_simps)
  also have "[... = h * (k - 1) * (1 * k + 1)] (mod 3)"
    by (intro cong_mult cong_refl cong_add cong_diff) (auto simp: Cong.cong_def)
  finally have cong: "[2 * dedekind_sum' h k = h * (k - 1) * (k + 1)] (mod
3)"
    by simp

  show "3 dvd (2 * dedekind_sum' h k) ⟷ ¬3 dvd k"
  proof (cases "3 dvd k")
    case True
    have "¬3 dvd h"
      using <coprime h k> True by fastforce
    have "[2 * dedekind_sum' h k = h * (k - 1) * (k + 1)] (mod 3)"
      by (fact cong)
    also have "[h * (k - 1) * (k + 1) = h * (0 - 1) * (0 + 1)] (mod 3)"
      using True by (intro cong_mult cong_diff cong_add cong_refl) (auto
simp: cong_0_iff)
    also have "h * (0 - 1) * (0 + 1) = -h"
      by simp
    finally have "3 dvd (2 * dedekind_sum' h k) ⟷ 3 dvd (-h)"

```

```

    using cong_dvd_iff by blast
  with <¬3 dvd h> True show ?thesis by auto
next
  case False
  hence "3 dvd (k + 1) ∨ 3 dvd (k - 1)"
    by presburger
  hence "3 dvd (h * (k - 1) * (k + 1))"
    by force
  also have "?this ↔ 3 dvd (2 * dedekind_sum' h k)"
    using cong cong_dvd_iff by blast
  finally show ?thesis
    using False by auto
qed
qed

lemma dedekind_sum'_reciprocity:
  fixes h k :: int
  assumes "h > 0" "k > 0" "coprime h k"
  shows "2 * h * dedekind_sum' h k = -2 * k * dedekind_sum' k h + h2
+ k2 - 3 * h * k + 1"
proof -
  have "real_of_int (2 * h * dedekind_sum' h k) = 12 * h * k * dedekind_sum
h k"
    unfolding of_int_mult of_int_dedekind_sum'[OF assms(2)] by simp
  also have "... = real_of_int (-12 * k * h) * dedekind_sum k h + real_of_int
(h2 + k2 - 3 * h * k + 1)"
    using dedekind_sum_reciprocity[OF assms] by simp
  also have "... = real_of_int (-2 * k * dedekind_sum' k h + h2 + k2 -
3 * h * k + 1)"
    using of_int_dedekind_sum'[OF assms(1)] by simp
  finally show ?thesis by linarith
qed

lemma cong_dedekind_sum'_1:
  fixes h k :: int
  defines "∅ ≡ gcd 3 k"
  assumes "h > 0" "coprime h k"
  shows "[2 * k * dedekind_sum' k h = 0] (mod ∅ * k)"
proof -
  have "∅ = 1 ∨ ∅ = 3"
    using gcd_prime_int[of 3 k] unfolding ∅_def by auto
  thus ?thesis
  proof
    assume "∅ = 3"
    hence "3 dvd k"
      unfolding ∅_def by (metis gcd_dvd2)
    with <coprime h k> have "¬3 dvd h"
      by force
  end

```

```

    have "3 dvd 2 * dedekind_sum' k h"
      using assms <¬3 dvd h>
      by (subst three_dvd_dedekind_sum'_iff_aux) (auto simp: coprime_commute)
    hence "3 * k dvd 2 * dedekind_sum' k h * k"
      by auto
    thus ?thesis
      using <∅ = 3> by (simp add: cong_0_iff)
  qed (auto simp: cong_0_iff)
qed

lemma cong_dedekind_sum'_2_aux:
  fixes h k :: int
  defines "∅ ≡ gcd 3 k"
  assumes "h > 0" "k > 0" "coprime h k"
  shows "[2 * h * dedekind_sum' h k = h2 + 1] (mod ∅ * k)"
proof -
  have "[-(2 * k * dedekind_sum' k h) + h2 + k2 - 3 * h * k + 1 = -0 +
h2 + 0 - 0 + 1] (mod ∅ * k)"
    unfolding ∅_def
    by (intro cong_diff cong_add cong_mult cong_refl cong_uminus cong_dedekind_sum'_1
  assms)
  (simp_all add: cong_0_iff power2_eq_square)
  also have "-(2 * k * dedekind_sum' k h) + h2 + k2 - 3 * h * k + 1 =
2 * h * dedekind_sum' h k"
    using dedekind_sum'_reciprocity[of k h] assms by (auto simp: coprime_commute)
  finally show ?thesis by simp
qed

lemma dedekind_sum'_negate:
  assumes "k > 0" "coprime h k"
  shows "dedekind_sum' (-h) k = -dedekind_sum' h k"
proof -
  have "real_of_int (dedekind_sum' (-h) k) = real_of_int (-dedekind_sum'
h k)"
    using assms unfolding of_int_minus
    by (subst (1 2) of_int_dedekind_sum') (auto simp: dedekind_sum_negate)
  thus ?thesis
    by linarith
qed

lemma cong_dedekind_sum'_2:
  fixes h k :: int
  defines "∅ ≡ gcd 3 k"
  assumes "k > 0" "coprime h k"
  shows "[2 * h * dedekind_sum' h k = h2 + 1] (mod ∅ * k)"
proof (cases h "0 :: int" rule: linorder_cases)
  case greater
  thus ?thesis
    using cong_dedekind_sum'_2_aux[of h k] assms by auto

```

```

next
  case less
  thus ?thesis
    using cong_dedekind_sum'_2_aux[of "-h" k] assms by (auto simp: dedekind_sum'_negate)
next
  case equal
  thus ?thesis using assms by (auto simp: Cong.cong_def)
qed

```

```

theorem dedekind_sum'_cong_8:
  assumes "k > 0" "coprime h k"
  shows "[2 * dedekind_sum' h k =
    (k-1)*(k+2) - 4*h*(k-1) + 4*( $\sum_{r \in \{1..<(k+1) \text{ div } 2\}} [2*h*r/k]$ )]
(mod 8)"
proof -
  define S1 where "S1 = ( $\sum_{r=1..<k} r * [of\_int (h * r) / k]$ )"
  define S2 where "S2 = ( $\sum_{r \mid r \in \{1..<k\} \wedge \text{odd } r} [of\_int (h * r) / k]$ )"
  define S3 where "S3 = ( $\sum_{r=1..<k} [of\_int (h * r) / k]$ )"
  define S4 where "S4 = ( $\sum_{r=1..<(k+1) \text{ div } 2} [of\_int (2 * h * r) / k]$ )"
  have "4 * 2 dvd 4 * (k * (k - 1))"
    by (intro mult_dvd_mono) auto
  hence dvd: "8 dvd 4 * k * (k - 1)"
    by (simp add: mult_ac)

  have "[4 * S1 = 4 * ( $\sum_{r=1..<k} \text{if even } r \text{ then } 0 \text{ else } [of\_int (h * r) / k]$ )] (mod 8)"
  unfolding S1_def sum_distrib_left
  proof (intro cong_sum, goal_cases)
    case (1 r)
    show ?case
    proof (cases "odd r")
      case True
      have *: "4 * r mod 8 = 4"
        using <odd r> by presburger
      have "[4 * r * [real_of_int (h * r) / k] = 4 * [real_of_int (h * r) / k]] (mod 8)"
        by (rule cong_mult[OF _ cong_refl]) (use * in <auto simp: Cong.cong_def>)
      thus ?thesis
        using True by (simp add: mult_ac)
    qed (auto simp: cong_0_iff)
  qed
  also have "( $\sum_{r=1..<k} \text{if even } r \text{ then } 0 \text{ else } [of\_int (h * r) / k]$ ) = S2"
  unfolding S2_def by (intro sum.mono_neutral_cong_right) auto
  finally have S12: "[4 * S1 = 4 * S2] (mod 8)" .

```

```

have "2 * dedekind_sum' h k =
      2 * (h * (k - 1) * (2 * k - 1)) - 12 * S1 - 6 * (k * (k - 1)
div 2)"
  using assms by (subst dedekind_sum'_altdef) (auto simp: S1_def)
  also have "6 * (k * (k - 1) div 2) = 3 * k * (k - 1)"
    by (subst div_mult_swap) auto
  also have "2 * (h * (k - 1) * (2 * k - 1)) - 12 * S1 - 3 * k * (k -
1) =
      -2 * h * (k - 1) + h * (4 * k * (k - 1)) - 12 * S1 + k *
(k - 1) - 4 * k * (k - 1)"
    by (simp add: algebra_simps)
  also have "[... = -2 * h * (k - 1) + h * 0 - 4 * S1 + k * (k - 1) - 0]
(mod 8)"
    using dvd by (intro cong_diff cong_add S12 cong_mult cong_refl)
      (auto simp: Cong.cong_def mod_eq_0_iff_dvd)
  also have "-2 * h * (k - 1) + h * 0 - 4 * S1 + k * (k - 1) - 0 = (k
- 1) * (k - 2 * h) - 4 * S1"
    by (simp add: algebra_simps)
  also have "[... = (k - 1) * (k - 2 * h) - 4 * S2] (mod 8)"
    by (intro cong_diff cong_refl S12)
  also have "S2 = S3 - S4"
  proof -
    have *: "{r. r ∈ {1..<k} ∧ odd r} = {1..<k} - {r. r ∈ {1..<k} ∧
even r}"
      by auto
    have "S2 = S3 - (∑ r | r ∈ {1..<k} ∧ even r. [of_int (h * r) / k])"
      unfolding S2_def S3_def * by (subst Groups_Big.sum_diff) auto
    also have "(∑ r | r ∈ {1..<k} ∧ even r. [of_int (h * r) / k]) = S4"
      unfolding S4_def using <k > 0>
      by (intro sum.reindex_bij_witness[of _ "λr. 2 * r" "λr. r div 2"])
      (auto simp: real_of_int_div)
    finally show ?thesis .
  qed
  also have "4 * (S3 - S4) = 4 * S3 - 4 * S4"
    by (simp add: algebra_simps)
  also have "4 * S3 = 2 * (h - 1) * (k - 1)"
  proof -
    have "real_of_int S3 = (∑ r=1..<k. -⟨of_int (h * r) / k⟩ + of_int
(h * r) / k - 1 / 2)"
      unfolding S3_def of_int_sum
    proof (intro sum.cong)
      fix r assume r: "r ∈ {1..<k}"
      hence "¬k dvd r"
        by (auto dest!: zdvd_imp_le)
      hence "¬k dvd (h * r)"
        using assms coprime_commute coprime_dvd_mult_right_iff by blast
      hence "real_of_int (h * r) / real_of_int k ∉ ℤ"
        using assms by (subst of_int_div_of_int_in_Ints_iff) auto
      thus "real_of_int [real_of_int (h * r) / real_of_int k] =

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      -(real_of_int (h * r) / real_of_int k) + real_of_int (h
* r) / real_of_int k - 1 / 2"
      by (simp add: dedekind_frac_def frac_def)
    qed auto
    also have "4 * ... = -4 * (∑ r=1..<k. ⟨of_int (h * r) / k⟩) +
      4 * h / k * (∑ r=1..<k. of_int r ^ 1) - ((real_of_int
k * 4 - 4) / 2)"
      using assms
      by (simp add: sum.distrib sum_subtractf sum_negf of_nat_diff mult_ac

      sum_distrib_left sum_distrib_right sum_divide_distrib)
    also have "(∑ r=1..<k. ⟨of_int (h * r) / k⟩) = 0"
      using assms by (intro sum_dedekind_frac_mult_eq_0)
    also have "4 * h / k * (∑ r=1..<k. real_of_int r ^ 1) = 2 * h * (k
- 1)"
      using assms by (subst sum_of_powers'_int_from_1) (auto simp: field_simps
bernpoly_def)
    also have "-4 * 0 + real_of_int (2 * h * (k - 1)) - (real_of_int k
* 4 - 4) / 2 =
      real_of_int (2 * (h - 1) * (k - 1))"
      by (simp add: field_simps)
    also have "4 * real_of_int S3 = real_of_int (4 * S3)"
      by simp
    finally show ?thesis by linarith
  qed
  also have "(k - 1) * (k - 2 * h) - (2 * (h - 1) * (k - 1) - 4 * S4)
=
      (k - 1) * (k + 2) - 4 * h * (k - 1) + 4 * S4"
      by (simp add: algebra_simps)
  finally show "[2 * dedekind_sum' h k = (k-1)*(k+2) - 4*h*(k-1) + 4*S4]
(mod 8)"
    unfolding S4_def .
qed

theorem dedekind_sum'_cong_8_odd:
  assumes "k > 0" "coprime h k" "odd k"
  shows "[2 * dedekind_sum' h k =
      k - 1 + 4*(∑ r∈{1..<(k+1) div 2}. [2*h*r/k])] (mod 8)"
proof -
  define S where "S = (∑ r=1..<(k+1) div 2. [of_int (2 * h * r) / k])"

  have 1: "[ (k-1)*(k+2) = k - 1 ] (mod 8)"
  proof -
    from assms obtain k' where k': "k = 2 * k' + 1"
      by (elim oddE)
    define k'' where "k'' = k' mod 4"
    have "k'' ∈ {0..<4}"
      unfolding k''_def by simp
    also have "{0..<4} = {0, 1, 2, 3::int}"

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```

    by auto
  finally have "(2 * k'' + 1) ^ 2 mod 8 = 1"
    by auto

  have "[k ^ 2 = ((2 * k') mod (2 * 4) + 1) ^ 2 mod 8] (mod 8)"
    unfolding k' by (intro cong_add cong_refl cong_pow cong_mod_rightI)
  (auto simp: Cong.cong_def)
  also have "(2 * k') mod (2 * 4) = 2 * k''"
    by (subst mod_mult_mult1) (auto simp: k''_def)
  also have "(2 * k'' + 1) ^ 2 mod 8 = 1"
    by fact
  finally have *: "[k ^ 2 = 1] (mod 8)" .

  have "(k-1)*(k+2) = k^2 + k - 2"
    by (simp add: algebra_simps power2_eq_square)
  also have "[... = 1 + k - 2] (mod 8)"
    by (intro cong_add cong_refl cong_diff *)
  finally show ?thesis
    by simp
qed

have 2: "[4 * h * (k - 1) = 0] (mod 8)"
proof -
  have "4 * 1 * 2 dvd 4 * h * (k - 1)"
    using assms by (intro mult_dvd_mono) auto
  thus ?thesis by (simp add: cong_0_iff)
qed

have "[2 * dedekind_sum' h k = (k-1)*(k+2) - 4*h*(k-1) + 4*S] (mod 8)"
  unfolding S_def using assms by (intro dedekind_sum'_cong_8)
also have "[ (k-1)*(k+2) - 4*h*(k-1) + 4*S = (k - 1) - 0 + 4 * S ] (mod 8)"
  by (intro cong_add cong_diff cong_refl 1 2)
finally show ?thesis
  by (simp add: S_def)
qed

lemma dedekind_sum'_cong_power_of_two:
  fixes h k k1 :: int and n :: nat
  assumes "h > 0" "k1 > 0" "odd k1" "n > 0" "k = 2 ^ n * k1" "coprime
h k"
  shows "[2 * h * dedekind_sum' h k =
          h^2 + k^2 + 1 + 5 * k - 4 * k * (∑ v=1..<(h+1) div 2. [of_int
(2 * k * v) / h])]
          (mod 2 ^ (n + 3))"
proof -
  from assms have "even k"

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    by auto
  with <coprime h k> have "odd h"
    using coprime_common_divisor odd_one by blast
  from assms have "k > 0"
    by auto
  define S where "S = (∑ v=1..<(h+1) div 2. [of_int (2 * k * v) / h])"
  have "[2 * dedekind_sum' k h * k = (h - 1 + 4 * S) * k] (mod (8 * k))"
    unfolding S_def using <h > 0> <k > 0> <coprime h k> <odd h>
    by (intro dedekind_sum'_cong_8_odd cong_cmult_rightI) (auto simp:
coprime_commute)
  also have "8 * k = 2 ^ (n + 3) * k1"
    using assms by (simp add: power_add)
  finally have *: "[2 * dedekind_sum' k h * k = (h - 1 + 4 * S) * k] (mod
2 ^ (n + 3))"
    using cong_modulus_mult by blast

  have **: "[k - 4 * h * k = 5 * k] (mod 2 ^ (n + 3))"
  proof -
    have "2 ^ 2 * 2 ^ n * 2 dvd 4 * k * (h + 1)"
      using assms by (intro mult_dvd_mono) auto
    hence "2 ^ (n + 3) dvd (5 * k - (k - 4 * h * k))"
      by (simp add: algebra_simps power_add)
    thus ?thesis
      by (subst cong_sym) (auto simp: cong_iff_dvd_diff)
  qed

  have "2 * h * dedekind_sum' h k = h2 + k2 - 3 * h * k + 1 - 2 * dedekind_sum'
k h * k"
    using dedekind_sum'_reciprocity[of h k] <h > 0> <k > 0> <coprime h
k> by simp
  also have "[... = h2 + k2 - 3 * h * k + 1 - (h - 1 + 4 * S) * k] (mod
2 ^ (n + 3))"
    using * by (intro cong_add cong_diff cong_refl)
  also have "h2 + k2 - 3*h*k + 1 - (h - 1 + 4*S) * k = h2 + k2 + 1 + (k
- 4*h*k) - 4*k*S"
    by (simp add: algebra_simps)
  also have "[... = h2 + k2 + 1 + 5 * k - 4*k*S] (mod 2 ^ (n + 3))"
    by (intro cong_add cong_diff cong_refl **)
  finally show ?thesis
    by (simp add: S_def)
  qed

lemma dedekind_sum'_cong_power_of_two':
  fixes h k k1 :: int
  assumes "h > 0" "k > 0" "even k" "coprime h k"
  shows " [2 * h * dedekind_sum' h k =
          h2 + k2 + 1 + 5 * k - 4 * k * (∑ v=1..<(h+1) div 2. [of_int
(2 * k * v) / h])]
          (mod 2 ^ (multiplicity 2 k + 3))"

```

```

proof (rule dedekind_sum'_cong_power_of_two)
  define k1 where "k1 = k div 2 ^ multiplicity 2 k"
  have "2 ^ multiplicity 2 k dvd k"
    using multiplicity_dvd by blast
  thus k_eq: "k = 2 ^ multiplicity 2 k * k1"
    by (auto simp: k1_def)
  show "odd k1"
    using <k > 0> multiplicity_decompose[of k 2] by (auto simp: k1_def)
  show "multiplicity 2 k > 0"
    using <even k > <k > 0> by (simp add: multiplicity_gt_zero_iff)
  show "k1 > 0"
    using <k > 0> by (subst (asm) k_eq) (auto simp: zero_less_mult_iff)
qed (use assms in auto)

```

```

lemma dedekind_sum_diff_even_int_aux:
  fixes a b c d :: int assumes det: "a * d - b * c = 1"
  fixes q c1 r δ' :: int and δ :: real
  assumes a: "a > 0"
  assumes q: "q ∈ {3, 5, 7, 13}" and "c1 > 0"
  assumes c: "c = q * c1"
  defines "r ≡ 24 div (q - 1)"
  defines "δ' ≡ (2 * dedekind_sum' a c - (a + d)) - (2 * q * dedekind_sum'
a c1 - (a + d) * q)"
  defines "δ ≡ (dedekind_sum a c - (a+d)/(12*c)) - (dedekind_sum a c1
- (a+d)/(12*c1))"
  shows "of_int δ' = 12 * c * δ" and "24 * c dvd r * δ'"
proof -
  define ϑ where "ϑ = gcd 3 c"
  define ϑ1 where "ϑ1 = gcd 3 c1"
  have "even r" "odd q"
    using q by (auto simp: r_def)

  have "q > 0"
    using q by auto
  have "c > 0"
    using q and <c1 > 0> and c by auto
  have "[a * d - b * 0 = a * d - b * c] (mod c)"
    by (intro cong_diff cong_mult cong_refl) (auto simp: Cong.cong_def)
  also have "a * d - b * c = 1"
    by fact
  finally have "[a * d = 1] (mod c)"
    by simp
  hence "coprime a c"
    using coprime_iff_invertible_int by blast
  have "coprime a c1"
    using <coprime a c> c by auto

```

```

show of_int_δ': "of_int δ' = 12 * c * δ"
  using <c > 0> <c1 > 0> <q > 0> unfolding δ'_def δ_def of_int_mult
of_int_diff
  by (subst (1 2) of_int_dedekind_sum') (auto simp: field_simps c)

have "[2 * a * dedekind_sum' a c = a2 + 1] (mod ϑ * c)"
  using <coprime a c> <c > 0> unfolding ϑ_def by (intro cong_dedekind_sum'_2)
auto
  hence "[2 * a * dedekind_sum' a c - a * (a + d) = (a2 + 1) - a * (a
+ d)] (mod ϑ * c)"
  by (intro cong_diff cong_refl)
  also have "(a2 + 1) - a * (a + d) = -b * c"
  using det unfolding power2_eq_square by (simp add: algebra_simps)
  finally have "[2 * a * dedekind_sum' a c - a * (a + d) = -b * c] (mod
ϑ * c)" .
  hence 1: "[2 * a * dedekind_sum' a c - a * (a + d) = -b * c] (mod ϑ1
* c)"
  by (rule cong_dvd_mono_modulus) (auto simp: ϑ_def ϑ1_def c)

have "[2 * a * dedekind_sum' a c1 * q = (a2 + 1) * q] (mod ϑ1 * c1
* q)"
  using <coprime a c1> <c1 > 0> unfolding ϑ1_def
  by (intro cong_cmult_rightI cong_dedekind_sum'_2) auto
  hence "[2 * a * dedekind_sum' a c1 * q - a * (a + d) * q =
(a2 + 1) * q - a * (a + d) * q] (mod ϑ1 * c)"
  by (intro cong_diff cong_refl) (auto simp: c mult_ac)
  also have "(a2 + 1) * q - a * (a + d) * q = -q * b * c"
  using det unfolding power2_eq_square by (simp add: algebra_simps)
  finally have 2: "[2 * a * q * dedekind_sum' a c1 - a * (a + d) * q =
-q * b * c] (mod ϑ1 * c)"
  by (simp add: mult_ac)

have r_δ': "r * δ' = r * (2 * dedekind_sum' a c - (a + d) -
(2 * q * dedekind_sum' a c1 - (a + d) * q))"
  by (simp add: δ'_def algebra_simps)

have r_a_δ': "r * a * δ' = r * (2 * a * dedekind_sum' a c - a * (a +
d) -
(2 * a * q * dedekind_sum' a c1 - a * (a + d) * q))"
  by (simp add: δ'_def algebra_simps)
  also have "[... = r * (-b * c - (-q * b * c))] (mod ϑ1 * c)"
  by (intro cong_mult cong_diff 1 2 cong_refl)
  also have "r * (-b * c - (-q * b * c)) = r * (q - 1) * b * c"
  by (simp add: algebra_simps)
  also have "r * (q - 1) = 24"
  using q by (auto simp: r_def)
  also have "ϑ1 * 1 * c dvd 24 * b * c"
  by (intro mult_dvd_mono) (auto simp: ϑ1_def gcd_dvdI1)
  hence "[24 * b * c = 0] (mod ϑ1 * c)"

```

```

    by (simp add: cong_0_iff)
  finally have "∅1 * c dvd r * δ' * a"
    by (simp add: cong_0_iff mult_ac)
  moreover have "coprime a (gcd 3 c1)"
    using <coprime a c1> coprime_imp_coprime dvd_trans by blast
  ultimately have "∅1 * c dvd r * δ'"
    using <coprime a c>
    by (subst (asm) coprime_dvd_mult_left_iff) (auto simp: ∅1_def coprime_commute)

  have "3 * c dvd r * δ'"
  proof (cases "q = 3")
    case False
    hence "¬3 dvd q"
      using q by auto
    hence "coprime 3 q"
      by (intro prime_imp_coprime) auto
    hence [simp]: "∅1 = ∅"
      by (auto simp: ∅_def ∅1_def c gcd_mult_right_left_cancel)
    show ?thesis
    proof (cases "∅ = 3")
      case True
      with <∅1 * c dvd r * δ'> show ?thesis by auto
    next
      case False
      hence [simp]: "∅ = 1"
        unfolding ∅_def by (subst gcd_prime_int) auto
      hence "¬3 dvd c"
        unfolding ∅_def by (subst (asm) gcd_prime_int) auto
      hence "¬3 dvd c1"
        unfolding c using <coprime 3 q> coprime_dvd_mult_right_iff by
blast
      have "[2 * dedekind_sum' a c = 0] (mod 3)"
        using three_dvd_dedekind_sum'_iff_aux[of c a] <c > 0> <coprime
a c> <¬3 dvd c>
        by (auto simp: cong_0_iff)
      moreover have "[2 * dedekind_sum' a c1 * q = 0] (mod 3)"
        using three_dvd_dedekind_sum'_iff_aux[of c1 a] <c1 > 0> <coprime
a c1> <¬3 dvd c1>
        by (auto simp: cong_0_iff)
      ultimately have "[r * δ' =
          r * ((0 - (a + d)) - (0 - (a + d) * q))] (mod
3)"
        unfolding r_δ' by (intro cong_mult cong_diff cong_refl) (simp_all
add: mult_ac)
      also have "r * ((0 - (a + d)) - (0 - (a + d) * q)) = r * (q - 1)
* (a + d)"
        by (simp add: algebra_simps)
      also have "r * (q - 1) = 24"
        using q by (auto simp: r_def)

```

```

also have "[24 * (a + d) = 0] (mod 3)"
  by (simp add: cong_0_iff)
finally have "3 dvd r * δ'"
  by (simp add: cong_0_iff)
moreover from ⟨∅1 * c dvd r * δ'⟩ have "c dvd r * δ'"
  by (simp add: mult_ac)
moreover have "coprime 3 c"
  using ⟨¬3 dvd c⟩ by (intro prime_imp_coprime) auto
ultimately show "3 * c dvd r * δ'"
  using divides_mult by blast
qed
next
case [simp]: True
have [simp]: "r = 12"
  by (simp add: r_def)
have [simp]: "∅ = 3"
  by (auto simp: ∅_def c)
show ?thesis
proof (cases "∅1 = 3")
  case True
  with ⟨∅1 * c dvd r * δ'⟩ show ?thesis by simp
next
  case False
  hence [simp]: "∅1 = 1"
    unfolding ∅1_def by (subst gcd_prime_int) auto
  hence "¬3 dvd c1"
    unfolding ∅1_def by (subst (asm) gcd_prime_int) auto
  have "[2 * dedekind_sum' a c1 * a * q = 0 * q] (mod 3 * q)"
    using three_dvd_dedekind_sum'_iff_aux[of c1 a] ⟨c1 > 0⟩ ⟨coprime
a c1⟩ ⟨¬3 dvd c1⟩
    by (intro cong_cmult_rightI) (auto simp: cong_0_iff)
  hence "[r * a * δ' = r * (2 * a * dedekind_sum' a c - a * (a + d)
- (0 - a * (a + d) * q))] (mod 9)"
    unfolding r_a_δ' by (intro cong_mult cong_diff cong_refl) (auto
simp: mult_ac)
  also have "r * (2 * a * dedekind_sum' a c - a * (a + d) - (0 - a
* (a + d) * q)) =
      r * (2 * a * dedekind_sum' a c) + 2 * r * (a2 + a * d)"
    by (simp add: algebra_simps power2_eq_square)
  also have "[2 * a * dedekind_sum' a c = a2 + 1] (mod ∅ * c)"
    using cong_dedekind_sum'_2[of c a] ⟨c > 0⟩ ⟨coprime a c⟩ unfold-
ing ∅_def by auto
  hence "[2 * a * dedekind_sum' a c = a2 + 1] (mod 9)"
    by (rule cong_dvd_mono_modulus) (auto simp: c)
  hence "[r * (2 * a * dedekind_sum' a c) + 2 * r * (a2 + a * d) =
r * (a2 + 1) + 2 * r * (a2 + a * d)] (mod 9)"
    by (intro cong_add cong_mult cong_refl)
  also have "r * (a2 + 1) + 2 * r * (a2 + a * d) = 3 * r * a2 + r
+ 2 * r * (a * d)"

```

```

    by (simp add: algebra_simps)
  also have "a * d = b * c + 1"
    using det by (simp add: algebra_simps)
  also have "3 * r * a2 + r + 2 * r * (b * c + 1) = 9 * (4 + 8 * b
* c1 + a2 * 4)"
    by (simp add: algebra_simps c)
  also have "[... = 0] (mod 9)"
    by (simp add: cong_0_iff)
  finally have "9 dvd r * a * δ'"
    by (simp add: cong_0_iff)
  moreover have "coprime a (3 ^ 2)"
    using <coprime a c> c by (subst coprime_power_right_iff) auto
  hence "coprime a 9"
    by (simp del: coprime_power_right_iff)
  ultimately have "9 dvd r * δ'"
    by (metis coprime_commute coprime_dvd_mult_right_iff mult.assoc
mult.commute)

  have "3 * c1 dvd 4 * δ'"
  proof (rule divides_mult)
    show "coprime 3 c1"
      using <ϑ1 = 1> unfolding ϑ1_def by auto
  next
    have "3 * c1 dvd 3 * (4 * δ)'"
      using <ϑ1 * c dvd r * δ'> by (simp add: c mult_ac)
    thus "c1 dvd 4 * δ'"
      by (subst (asm) dvd_mult_cancel_left) auto
  next
    have "3 * 3 dvd 3 * (4 * δ)'"
      using <9 dvd r * δ'> by simp
    thus "3 dvd 4 * δ'"
      by (subst (asm) dvd_mult_cancel_left) auto
  qed
  hence "3 * c dvd 3 * (4 * δ)'"
    by (intro mult_dvd_mono dvd_refl) (auto simp: c)
  thus ?thesis
    by (simp add: c mult_ac)
  qed
qed

show "24 * c dvd r * δ'"
proof (cases "even c")
  assume "odd c"
  hence "odd c1"
    by (auto simp: c)

  define T :: "int ⇒ int" where
    "T = (λc. (∑ r = 1..<(c + 1) div 2. [real_of_int (2 * a * r) / real_of_int
c])))"

```

```

have "[2 * dedekind_sum' a c = c - 1 + 4 * T c] (mod 8)"
  unfolding T_def by (rule dedekind_sum'_cong_8_odd)
  (use <coprime a c> <odd c> <c > 0> in auto)
moreover have "[2 * dedekind_sum' a c1 * q = (c1 - 1 + 4 * T c1)
* q] (mod 8)"
  unfolding T_def by (intro cong_mult cong_refl dedekind_sum'_cong_8_odd)
  (use <coprime a c1> <odd c1> <c1 > 0> in auto)
ultimately have "[r * δ' = r * ((c - 1 + 4 * T c - (a + d)) -
((c1 - 1 + 4 * T c1) * q - (a + d)
* q))] (mod 8)"
  unfolding δ'_def by (intro cong_diff cong_mult[of r] cong_refl)
(auto simp: mult_ac)
also have "r * ((c - 1 + 4*T c - (a+d)) - ((c1 - 1 + 4*T c1)*q - (a+d)*q))
=
      r * (q - 1) * (a + d + 1) + (r * 4) * (T c - q * T c1)"
  by (simp add: c algebra_simps)
also have "r * (q - 1) = 24"
  using q by (auto simp: r_def)
also have "[24 * (a + d + 1) + r * 4 * (T c - q * T c1) =
0 * (a + d + 1) + 0 * (T c - q * T c1)] (mod 8)"
  using <even r> by (intro cong_add cong_mult cong_refl) (auto simp:
cong_0_iff)
finally have "8 dvd r * δ'"
  by (simp add: cong_0_iff)

have "8 * (3 * c) dvd r * δ'"
proof (rule divides_mult)
  have "coprime (2 ^ 3) (3 * c)"
    using <odd c> unfolding coprime_power_left_iff by auto
  thus "coprime 8 (3 * c)"
    by (simp del: coprime_power_left_iff)
qed fact+
thus ?thesis
  by simp
next
assume "even c"
with <coprime a c> have "odd a"
  using coprime_common_divisor odd_one by blast
from <even c> and <odd q> have "even c1"
  by (auto simp: c)

define n where "n = multiplicity 2 c"
define c' where "c' = c div 2 ^ n"
have "c = 2 ^ n * c'"
  unfolding c'_def n_def by (simp add: multiplicity_dvd)
have n_altdef: "n = multiplicity 2 c1"
  using <odd q> by (auto simp: n_def c multiplicity_prime_elem_times_other)

```

```

have "odd c'"
  unfolding c'_def n_def using <c > 0> multiplicity_decompose[of c
2] by auto

define T where "T = (λc. (∑ v = 1..<(a + 1) div 2. [real_of_int (2
* c * v) / real_of_int a]))"

have "[2 * a * dedekind_sum' a c = a2 + c2 + 1 + 5 * c - 4 * c * T
c] (mod 2 ^ (n + 3))"
  unfolding n_def T_def using <a > 0> <c > 0> <coprime a c> <even
c> <odd a>
  by (intro dedekind_sum'_cong_power_of_two') auto
  moreover have "[2 * a * dedekind_sum' a c1 * q = (a2 + c12 + 1 +
5 * c1 - 4 * c1 * T c1) * q]
(mod 2 ^ (n + 3))"
  unfolding n_altdef T_def using <a > 0> <c1 > 0> <coprime a c1>
<even c1> <odd a>
  by (intro cong_mult[of _ _ _ q] dedekind_sum'_cong_power_of_two')
auto
ultimately have "[r * a * δ' =
r * ((a2 + c2 + 1 + 5 * c - 4 * c * T c) - a *
(a + d) -
((a2 + c12 + 1 + 5 * c1 - 4 * c1 * T c1) * q -
a * (a + d) * q))]
(mod 2 ^ (n + 3))"
  unfolding r_a_δ' by (intro cong_mult[of r] cong_diff cong_refl)
(auto simp: mult_ac)
  also have "r * ((a2 + c2 + 1 + 5 * c - 4 * c * T c) - a * (a + d)
-
((a2 + c12 + 1 + 5 * c1 - 4 * c1 * T c1) * q -
a * (a + d) * q)) =
r*(q-1) * (a * d - 1 + c * c1) - 4*c*r * (T c - T c1)"
  by (simp add: algebra_simps c power2_eq_square)
  also have "r * (q - 1) = 24"
  using q by (auto simp: r_def)
  also have "a * d - 1 = b * c"
  using det by (simp add: algebra_simps)
  also have "24 * (b * c + c * c1) = 24 * c * (b + c1)"
  by (simp add: algebra_simps)
  also have "[24 * c * (b + c1) - 4 * c * r * (T c - T c1) = 0 - 0]
(mod 2 ^ (n + 3))"
  proof (intro cong_diff)
    have "2 ^ (n + 3) dvd 2 ^ (n + 3) * (3 * c' * (b + c1))"
    using dvd_triv_left by blast
    also have "... = 24 * c * (b + c1)"
    by (simp add: <c = 2 ^ n * c'> mult_ac power_add)
    finally show "[24 * c * (b + c1) = 0] (mod 2 ^ (n + 3))"
    by (simp add: cong_0_iff)
  next

```

```

      have "4 * 2 ^ n * 2 * 1 dvd 4 * c * r * (T c - T c1)"
        using <even r> by (intro mult_dvd_mono) (auto simp: <c = 2 ^ n
* c'>)
      thus "[4 * c * r * (T c - T c1) = 0] (mod 2 ^ (n + 3))"
        by (simp add: power_add cong_0_iff mult_ac)
    qed
    finally have "2 ^ (n + 3) dvd r * δ' * a"
      by (simp add: cong_0_iff mult_ac)
    hence "2 ^ (n + 3) dvd r * δ'"
      using <odd a> by (subst (asm) coprime_dvd_mult_left_iff) auto

    have "2 ^ (n + 3) * (3 * c') dvd r * δ'"
    proof (rule divides_mult)
      show "coprime (2 ^ (n + 3)) (3 * c')"
        using <odd c'> by simp
      show "2 ^ (n + 3) dvd r * δ'"
        by fact
      have "3 * c' dvd 3 * c"
        by (auto simp: <c = 2 ^ n * c'>)
      also have "3 * c dvd r * δ'"
        by fact
      finally show "3 * c' dvd r * δ'" .
    qed
    thus ?thesis
      by (simp add: <c = 2 ^ n * c'> power_add mult_ac)
  qed
qed

```

```

theorem dedekind_sum_diff_even_int:
  fixes a b c d :: int assumes det: "a * d - b * c = 1"
  fixes q c1 r :: int and δ' :: "int ⇒ int" and δ :: "int ⇒ real"
  assumes q: "q ∈ {3, 5, 7, 13}" and "c1 > 0"
  assumes c: "c = q * c1"
  defines "r ≡ 24 div (q - 1)"
  defines "δ' ≡ (λa. 2 * dedekind_sum' a c - (a + d) - (2 * q * dedekind_sum'
a c1 - (a + d) * q))"
  defines "δ ≡ (λa. dedekind_sum a c - (a+d)/(12*c) - (dedekind_sum a
c1 - (a+d)/(12*c1)))"
  shows "of_int (δ' a) = 12 * c * δ a"
    and "24 * c dvd r * δ' a"
    and "real_of_int r * δ a / 2 ∈ ℤ"
proof -
  define a' t where "a' = a mod c" and "t = a div c"
  have a'_eq: "a' = a - t * c"
    by (simp add: a'_def t_def algebra_simps)
  have cong1: "[a' = a] (mod c)"
    by (simp add: a'_def)
  hence cong2: "[a' = a] (mod c1)"

```

```

using c by (metis cong_modulus_mult mult.commute)

have "q > 0"
  using q by auto
have "c > 0"
  using q and <c1 > 0> and c by auto
have "[a * d - b * 0 = a * d - b * c] (mod c)"
  by (intro cong_diff cong_mult cong_refl) (auto simp: Cong.cong_def)
also have "a * d - b * c = 1"
  by fact
finally have "[a * d = 1] (mod c)"
  by simp
hence "coprime a c"
  using coprime_iff_invertible_int by blast
have "coprime a c1"
  using <coprime a c> c by auto

have "a' ≥ 0"
  using <c > 0> by (auto simp: a'_def)
moreover have "a' ≠ 0"
  using <coprime a c> q by (auto simp: a'_def c)
ultimately have "a' > 0"
  by linarith

have 1: "dedekind_sum a' c = dedekind_sum a c" using cong1
  by (rule dedekind_sum_cong) (use <coprime a c> in <auto simp: a'_def
coprime_commute>)
have 2: "dedekind_sum' a' c = dedekind_sum' a c" using cong1
  by (rule dedekind_sum'_cong) (use <coprime a c> in <auto simp: a'_def
coprime_commute>)
have 3: "dedekind_sum a' c1 = dedekind_sum a c1" using cong2
  by (rule dedekind_sum_cong) (use <coprime a c1> in <auto simp: a'_def
coprime_commute>)
have 4: "dedekind_sum' a' c1 = dedekind_sum' a c1" using cong2
  by (rule dedekind_sum'_cong) (use <coprime a c1> in <auto simp: a'_def
coprime_commute>)

have δ_eq: "δ a' = δ a - t * (q - 1) / 12"
  unfolding δ_def 1 3 using <c > 0> <c1 > 0> <q > 0>
  by (simp add: field_simps c a'_eq)
have δ'_eq: "δ' a' = δ' a - c * t * (q - 1)"
  unfolding δ'_def 2 4 using <c > 0> <c1 > 0> <q > 0>
  by (simp add: field_simps c a'_eq)

have det': "a' * d - (b - t * d) * c = 1"
  using det by (simp add: a'_eq algebra_simps)

have "of_int (δ' a') = 12 * c * δ a'"
  unfolding δ'_def δ_def using det' <c1 > 0> q c <a' > 0>

```

```

    by (intro dedekind_sum_diff_even_int_aux[of a' d "b - t * d" c]) auto
  thus of_int_δ': "of_int (δ' a) = 12 * c * δ a"
    by (simp add: δ_eq δ'_eq field_simps)

  have "24 * c dvd r * δ' a"
    unfolding r_def δ'_def using det' <c1 > 0> q c <a' > 0>
    by (intro dedekind_sum_diff_even_int_aux[of a' d "b - t * d" c]) auto
  also have "r * δ' a = r * δ' a - (q - 1) * r * c * t"
    by (simp add: δ'_eq algebra_simps)
  also have "(q - 1) * r = 24"
    unfolding r_def using q by auto
  also have "24 * c dvd r * δ' a - 24 * c * t ⟷ 24 * c dvd r * δ' a"
    by (rule dvd_diff_left_iff) auto
  finally show "24 * c dvd r * δ' a" .

  then obtain m where m: "r * δ' a = 24 * c * m"
    by auto
  hence "real_of_int (r * δ' a) = real_of_int (24 * c * m)"
    by (simp only: )
  hence "real_of_int r * δ a / 2 = of_int m"
    using <c > 0> by (simp add: of_int_δ' field_simps)
  also have "... ∈ ℤ"
    by simp
  finally show "real_of_int r * δ a / 2 ∈ ℤ" .
qed

no_notation dedekind_frac ("⟨_⟩")

end

```

References

- [1] T. M. Apostol. *Modular Functions and Dirichlet Series in Number Theory*. Graduate Texts in Mathematics. Springer, 1990.