

Coppersmith’s Method

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Abstract

We formalize *Coppersmith’s method*, an algorithm for finding small (in magnitude) roots of univariate integer polynomials mod M . Coppersmith’s method has important applications in cryptography and is used in various attacks on the RSA algorithm for public-key cryptography. We also formalize a related (more lightweight) result with slightly weaker bounds; we split out the generic mathematical results underlying both this lightweight result and Coppersmith’s method into a dedicated locale, which could be used to prove other “Coppersmith-like” results. Our work builds on the existing formalization of the LenstraLenstraLovász (LLL) algorithm for lattice basis reduction [2], and our formalization adds a determinant bound on the length of the short vector produced by the LLL algorithm.

There are many resources on Coppersmith’s method that we found useful in the course of our development. Some that we recommend include Chapter 19 of a textbook by Galbraith [3], Section 17.3 of a textbook by Trappe and Washington [4], and the following set of lecture notes [1].

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1 Preliminary results for LLL

In this file, we prove some additional results involving lattices and a bound for LLL.

```
theory More-LLL
```

```
imports
```

```
  LLL-Basis-Reduction.LLL
```

```
begin
```

```
context vec-module begin
```

```
lemma in-lattice-ofD:
```

```
  assumes x:  $x \in \text{lattice-of } a$ 
```

```
  assumes a:  $\text{set } a \subseteq \text{carrier-vec } n$ 
```

```
  obtains v where
```

```
     $v \in \text{carrier-vec } (\text{length } a)$ 
```

```
     $\text{mat-of-cols } n \ a \ *_{\mathbf{v}} \ \text{map-vec of-int } (v::\text{int vec}) = x$ 
```

```
  <proof>
```

lemma *exists-list-all2*:

assumes $\forall x. (x \in \text{set } xs \longrightarrow (\exists y. P \ x \ y))$

obtains *ys* **where** *list-all2* *P* *xs* *ys*

<proof>

lemma *subset-lattice-ofD*:

assumes *xs*: *set* *xs* \subseteq *lattice-of a set* *xs* \subseteq *carrier-vec* *n*

assumes *a*: *set* *a* \subseteq *carrier-vec* *n*

obtains *vs* **where**

vs \in *carrier-mat* (*length* *a*) (*length* *xs*)

mat-of-cols *n* *a* * *map-mat of-int* *vs* = *mat-of-cols* *n* *xs*

<proof>

lemma *mk-coeff-list-nth*:

assumes *i* < *length* *ls* *distinct* *ls*

shows *mk-coeff* *ls* *f* (*ls* ! *i*) = *f* *i*

<proof>

This next lemma is trivial when the cols are not distinct.

lemma *mat-mul-non-zero-col-lin-dep*:

assumes *A*: *A* \in *carrier-mat* *n* *y* *distinct* (*cols* *A*)

assumes *U*: *U* \in *carrier-mat* *y* *z*

assumes *i*: *i* < *z* *col* *U* *i* \neq 0_v *y*

assumes *eqz*: *A* * *U* = 0_m *n* *z*

shows *lin-dep* (*set* (*cols* *A*))

<proof>

lemma *lin-indpt-mul-eq-ident*:

assumes *a*: *distinct* *a* *lin-indpt* (*set* *a*)

set *a* \subseteq *carrier-vec* *n* *length* *a* = *m*

assumes *u*: *u* \in *carrier-mat* *m* *m*

assumes *e*: *mat-of-cols* *n* *a* = *mat-of-cols* *n* *a* * *u*

shows *u* = 1_m *m*

<proof>

end

Set up a locale: *vec_module* where the ring has characteristic zero

locale *ring-char-0-vec-module* = *vec-module* *f-ty* *n* **for**

f-ty::'*a*:: {*comm-ring-1*, *ring-char-0*} *itself*

and *n*

begin

This next lemma shows that different basis *a*, *b* of the same lattice have the same Gram determinant.

lemma *lattice-of-eq-gram-det-eq*:

fixes *a* *b*::'*a* *vec* *list*

assumes *a*: *distinct* *a* *lin-indpt* (*set* *a*) *set* *a* \subseteq *carrier-vec* *n*

assumes *b*: *set* *b* \subseteq *carrier-vec* *n* *length* *a* = *length* *b*

assumes *lat-eq*: *lattice-of a = lattice-of b*
defines *A*: $A \equiv \text{mat-of-cols } n \ a$
defines *B*: $B \equiv \text{mat-of-cols } n \ b$
shows $\det (A^T * A) = \det (B^T * B)$
 <proof>

lemma *lattice-of-eq-gram-det-rows-eq*:
fixes *a b*: 'a vec list
assumes *a*: *distinct a lin-indpt (set a) set a ⊆ carrier-vec n*
assumes *b*: *set b ⊆ carrier-vec n length a = length b*
assumes *lat-eq*: *lattice-of a = lattice-of b*
defines *A*: $A \equiv \text{mat-of-rows } n \ a$
defines *B*: $B \equiv \text{mat-of-rows } n \ b$
shows $\det (A * A^T) = \det (B * B^T)$
 <proof>

lemma *lattice-of-eq-sq-det-eq*:
fixes *a b*: 'a vec list
assumes *a*: *distinct a lin-indpt (set a) set a ⊆ carrier-vec n length a = n*
assumes *b*: *set b ⊆ carrier-vec n length b = n*
assumes *lat-eq*: *lattice-of a = lattice-of b*
shows $(\det (\text{mat-of-cols } n \ a))^2 = (\det (\text{mat-of-cols } n \ b))^2$
 <proof>

lemma *lattice-of-eq-sq-det-rows-eq*:
fixes *a b*: 'a vec list
assumes *a*: *distinct a lin-indpt (set a) set a ⊆ carrier-vec n length a = n*
assumes *b*: *set b ⊆ carrier-vec n length b = n*
assumes *lat-eq*: *lattice-of a = lattice-of b*
shows $(\det (\text{mat-of-rows } n \ a))^2 = (\det (\text{mat-of-rows } n \ b))^2$
 <proof>

end

context *LLL-with-assms*
begin

This next lemma bounds the size of the shortest vector by the determinant.

lemma *short-vector-det-bound*:
assumes *m*: $m \neq 0$
assumes *k*: $k \leq m$
shows
 $(\text{rat-of-int } \|\text{short-vector}\|^2) \wedge k \leq$
 $\alpha \wedge (k * (k-1) \text{ div } 2) * \text{rat-of-int } (\text{gs.Gramian-determinant reduce-basis } k)$
 <proof>

lemma *square-Gramian-determinant-eq-det-square*:
assumes *sq:n = m*

```

shows gs.Gramian-determinant fs-init m =
  (det (mat-of-rows m fs-init))2
  ⟨proof⟩

```

```

end

```

```

end

```

2 Algorithm for Coppersmith's Method

In this file, we formalize the algorithm for Coppersmith's method. We follow the descriptions in Chapter 19 of "Mathematics of Public Key Cryptography" by Steven Galbraith and Section 17.3 of "Introduction to Cryptography with Coding Theory" by Trappe and Washington. We first formalize an algorithm for a "lightweight" version of Coppersmith's method, and then formalize Coppersmith's method itself. Correctness proofs for these algorithms are in "Towards_Coppersmith.thy" and "Coppersmith.thy".

```

theory Coppersmith-Algorithm

```

```

imports LLL-Factorization.LLL-Factorization
begin

```

This next definition forms the vector associated to a polynomial as in (19.1) of Galbraith.

```

definition vec-associated-to-poly:: ('a::{ring,power}) poly ⇒ 'a ⇒ 'a vec where
  vec-associated-to-poly p X = vec (degree p + 1) (λj. (coeff p j) * Xj)

```

```

abbreviation vec-associated-to-int-poly:: int poly ⇒ int ⇒ int vec where
  vec-associated-to-int-poly ≡ vec-associated-to-poly

```

```

abbreviation vec-associated-to-real-poly:: real poly ⇒ real ⇒ real vec where
  vec-associated-to-real-poly ≡ vec-associated-to-poly

```

```

definition int-poly-associated-to-vec:: int vec ⇒ nat ⇒ int poly where
  int-poly-associated-to-vec v X = (
    let newvec = vec (dim-vec v) (λj. floor (((v $ j)::real)/(Xj))) in
    ∑ i<(dim-vec v). monom (newvec $ i) i
  )

```

2.1 Lightweight method similar to Coppersmith

In this section, we start with a more lightweight version of Coppersmith which doesn't achieve the full bounds, but is built on similar ideas.

This next definition constructs the g_i 's

```

definition g-i-vec:: nat ⇒ nat ⇒ nat ⇒ nat ⇒ int vec where
  g-i-vec M X i n = vec n (λj. if j = i then M*Xi else 0)

```

This next definition should be called with degree $d = p - 1$

definition *form-basis-helper*:: $int\ poly \Rightarrow nat \Rightarrow nat \Rightarrow nat \Rightarrow int\ vec\ list$ **where**
form-basis-helper $p\ M\ X\ d = map\ (\lambda i. g-i-vec\ M\ X\ i\ (degree\ p + 1))\ [0..<d + 1]$

definition *form-basis*:: $int\ poly \Rightarrow nat \Rightarrow nat \Rightarrow nat \Rightarrow int\ vec\ list$ **where**
form-basis $p\ M\ X\ d = (form-basis-helper\ p\ M\ X\ d)\ @\ [vec-associated-to-int-poly\ p\ X]$

definition *first-vector*:: $int\ poly \Rightarrow nat \Rightarrow nat \Rightarrow int\ vec$ **where**
first-vector $p\ M\ X = (\$
 $\quad let\ cs-basis = form-basis\ p\ M\ X\ (degree\ p - 1)\ in$
 $\quad (short-vector\ 2\ cs-basis)$
 $\quad)$

definition *towards-coppersmith*:: $int\ poly \Rightarrow nat \Rightarrow nat \Rightarrow int\ poly$ **where**
towards-coppersmith $p\ M\ X = int-poly-associated-to-vec\ (first-vector\ p\ M\ X)\ X$

2.2 Full Coppersmith

In this section, we develop the full Coppersmith's algorithm.

definition *vec-associated-to-int-poly-padded*:: $nat \Rightarrow int\ poly \Rightarrow nat \Rightarrow int\ vec$ **where**
vec-associated-to-int-poly-padded $n\ p\ X = vec\ n\ (\lambda j. (coeff\ p\ j)*X^j)$

definition *row-of-coppersmith-matrix*:: $int\ poly \Rightarrow nat \Rightarrow nat \Rightarrow nat \Rightarrow nat \Rightarrow int\ vec$
 $\Rightarrow int\ vec$
where *row-of-coppersmith-matrix* $p\ M\ X\ h\ i\ j =$
vec-associated-to-int-poly-padded $((degree\ p)*h)\ (smult\ (((M^{h-1-j})))\ (p^j*(monom\ 1\ i)))\ X$

definition *form-basis-coppersmith-aux*:: $int\ poly \Rightarrow nat \Rightarrow nat \Rightarrow nat \Rightarrow int\ vec\ list$
where *form-basis-coppersmith-aux* $p\ M\ X\ h\ i\ j = (map\ (\lambda i. row-of-coppersmith-matrix\ p\ M\ X\ h\ i\ j)\ [0..<degree\ p])$

definition *form-basis-coppersmith*:: $int\ poly \Rightarrow nat \Rightarrow nat \Rightarrow int\ vec\ list$
where *form-basis-coppersmith* $p\ M\ X\ h = concat\ (map\ (\lambda j. form-basis-coppersmith-aux\ p\ M\ X\ h\ j)\ [0..<(h::nat)])$

definition *calculate-h-coppersmith-aux*:: $int\ poly \Rightarrow real \Rightarrow int$ **where**
calculate-h-coppersmith-aux $p\ e = (let\ d = degree\ p\ in\ ((ceiling\ (((d-1)/(d*(e::real))\ + 1)/d))::int))$

definition *calculate-h-coppersmith*:: $int\ poly \Rightarrow real \Rightarrow nat$ **where**
calculate-h-coppersmith $p\ e = (nat\ (calculate-h-coppersmith-aux\ p\ e))$

Note that we pass 2 as a parameter to the LLL algorithm. Any bound $> 4/3$ would work.

definition *first-vector-coppersmith*:: $int\ poly \Rightarrow nat \Rightarrow nat \Rightarrow real \Rightarrow int\ vec$ **where**
first-vector-coppersmith $p\ M\ X\ epsilon = (\$

$let\ cs\text{-}basis = form\text{-}basis\text{-}coppersmith\ p\ M\ X\ (calculate\text{-}h\text{-}coppersmith\ p\ epsilon)$
in
 $short\text{-}vector\ 2\ cs\text{-}basis)$

definition $coppersmith::\ int\ poly \Rightarrow nat \Rightarrow nat \Rightarrow real \Rightarrow int\ poly$ **where**
 $coppersmith\ p\ M\ X\ epsilon =$
 $int\text{-}poly\text{-}associated\text{-}to\text{-}vec\ (first\text{-}vector\text{-}coppersmith\ p\ M\ X\ epsilon)\ X$

end

3 Howgrave-Graham's theorem

In this file, we prove a result due to Howgrave-Graham on small-enough roots of polynomials mod M (see Theorem 19.1.2 in "Mathematics of Public Key Cryptography" by Galbraith).

theory *Howgrave-Graham*

imports *Coppersmith-Algorithm*
HOL-Analysis.L2-Norm
LLL-Basis-Reduction.Norms

begin

abbreviation $euclidean\text{-}norm\text{-}int\text{-}vec::int\ vec \Rightarrow real$
where $euclidean\text{-}norm\text{-}int\text{-}vec\ v \equiv sqrt\ (sq\text{-}norm\text{-}vec\ v)$

abbreviation $euclidean\text{-}norm\text{-}real\text{-}vec::real\ vec \Rightarrow real$
where $euclidean\text{-}norm\text{-}real\text{-}vec\ v \equiv sqrt\ (sq\text{-}norm\text{-}vec\ v)$

lemma $euclidean\text{-}norm\text{-}int\text{-}vec\text{-}eq:$
shows $euclidean\text{-}norm\text{-}int\text{-}vec\ v = sqrt\ (\sum\ i<(dim\text{-}vec\ v). (v\$i)^2)$
 $\langle proof \rangle$

lemma $euclidean\text{-}norm\text{-}real\text{-}vec\text{-}eq:$
shows $sqrt\ (sq\text{-}norm\text{-}vec\ v) = sqrt\ (\sum\ i<(dim\text{-}vec\ v). (v\$i)^2)$
 $\langle proof \rangle$

lemma $euclidean\text{-}norm\text{-}gteq0:$
shows $euclidean\text{-}norm\text{-}real\text{-}vec\ (a::real\ vec) \geq 0$
 $euclidean\text{-}norm\text{-}int\text{-}vec\ (c::int\ vec) \geq 0$
 $\langle proof \rangle$

lemma $dim\text{-}vec\text{-}vec\text{-}associated\text{-}to\text{-}poly[simp]:$
shows $dim\text{-}vec\ (vec\text{-}associated\text{-}to\text{-}poly\ F\ X) = degree\ F + 1$
 $\langle proof \rangle$

lemma $Cauchy\text{-}Schwarz\text{-}sum:$
fixes $n::\ nat$
fixes $x::\ nat \Rightarrow real$

shows $(\sum_{i \leq n}. x\ i) \leq \text{sqrt} ((n+1) * (\sum_{i \leq n}. (x\ i)^2))$
 <proof>

lemma *abs-mult-sum*:
fixes $f\ g :: \text{nat} \Rightarrow \text{real}$
fixes $n :: \text{nat}$
shows $\text{abs}(\sum_{i \leq n}. (f\ i)*(g\ i))$
 $\leq (\sum_{i \leq n}. (\text{abs}\ (f\ i))*(\text{abs}\ (g\ i)))$
 <proof>

lemma *sum-helper*:
fixes $g\ h :: \text{nat} \Rightarrow \text{real}$
assumes $\forall i \leq n. f\ i \geq 0$
assumes $\forall i \leq n. g\ i \leq h\ i$
shows $(\sum_{i \leq n}. (f\ i)*(g\ i)) \leq (\sum_{i \leq n}. (f\ i)*(h\ i))$
 <proof>

theorem *Howgrave-Graham*:
fixes $F :: \text{real poly}$
fixes $M\ X :: \text{nat}$
fixes $x0\ k :: \text{int}$
assumes *M-gt*: $M > 0$
assumes *root-mod-M*: $\text{poly}\ F\ (\text{real-of-int}\ x0) = k * M$
assumes *root-bound*: $\text{abs}\ x0 \leq X$
assumes *norm-bound*: $\text{sqrt} (\|\text{vec-associated-to-poly}\ F\ X\|^2) < M / \text{sqrt} (\text{degree}\ F + 1)$
shows $\text{poly}\ F\ x0 = 0$
 <proof>

abbreviation *int-poly-to-real-poly*:: $\text{int poly} \Rightarrow \text{real poly}$
where *int-poly-to-real-poly* $F \equiv \text{map-poly}\ \text{real-of-int}\ F$

lemma *int-poly-to-real-poly-same-norm*:
fixes $X :: \text{nat}$
shows $\text{euclidean-norm-int-vec}\ (\text{vec-associated-to-int-poly}\ F\ X) =$
 $\text{euclidean-norm-real-vec}\ (\text{vec-associated-to-real-poly}\ (\text{int-poly-to-real-poly}\ F)\ X)$
 <proof>

Now we restate the result over int polys.

lemma *Howgrave-Graham-int-poly*:
fixes $F :: \text{int poly}$
fixes $M\ X :: \text{nat}$
fixes $x0 :: \text{int}$
assumes *M-gt*: $M > 0$
assumes *root-mod-M*: $\text{poly}\ F\ x0 \text{ mod } M = 0$
assumes *root-bound*: $\text{abs}\ x0 \leq X$
assumes *norm-bound*: $\text{sqrt} (\text{sq-norm-vec}\ (\text{vec-associated-to-int-poly}\ F\ X)) < M / \text{sqrt} (\text{degree}\ F + 1)$

shows $\text{poly } F \ x0 = 0$
 ⟨*proof*⟩

end

4 Coppersmith Generic

In this file, we develop the generic argument behind Coppersmith's method.

theory *Coppersmith-Generic*

imports *Coppersmith-Algorithm*

Howgrave-Graham

More-LLL

begin

hide-const (open) *module.smult*

4.1 Some matrix properties

This definition should only be used for lists of vectors that correspond to square matrices

definition *vec-list-to-square-mat*:: 'a vec list \Rightarrow 'a mat
where *vec-list-to-square-mat* L = *mat-of-rows* (length L) L

This lemma is similar to `upper_triangular_imp_distinct`, from the "Jordan_Normal_Form.Matrix" library by Thiemann and Yamada.

lemma *lower-triangular-imp-distinct*:

assumes A: A \in *carrier-mat* n n

and *tri*: $\bigwedge i j. i < j \Rightarrow j < n \Rightarrow A \ \S\ (i,j) = 0$

and *diag*: $0 \notin \text{set}(\text{diag-mat } A)$

shows *distinct* (rows A)

⟨*proof*⟩

lemma *lower-triangular-imp-det-eq-0-iff*:

fixes A :: 'a :: *idom* mat

assumes A \in *carrier-mat* n n **and** $\bigwedge i j. i < j \Rightarrow j < n \Rightarrow A \ \S\ (i,j) = 0$

shows $\det A = 0 \iff 0 \in \text{set}(\text{diag-mat } A)$

⟨*proof*⟩

This next lemma is similar to `upper_triangular_imp_lin_indpt_rows`, from the "Jordan_Normal_Form.VS_Connect" library by Thiemann and Yamada.

lemma (in *idom-vec*) *lower-triangular-imp-lin-indpt-rows*:

assumes A: A \in *carrier-mat* n n

and *lower-tri*: $\bigwedge i j. i < j \Rightarrow j < n \Rightarrow A \ \S\ (i,j) = 0$

and *diag*: $0 \notin \text{set}(\text{diag-mat } A)$

shows *lin-indpt* (set (rows A))

<proof>

lemma *g-i-vec-ith-element:*

assumes *degree* $p \geq 1$

assumes $i < \text{degree } p$

assumes $M > 0$

assumes $X > 0$

shows $(g\text{-}i\text{-}vec\ M\ X\ i\ (\text{degree } p + 1))\ \$\ i = M * X^{\wedge} i$

<proof>

lemma *ith-row-form-basis-helper:*

assumes $i \leq d$

shows $((\text{form-basis-helper } p\ M\ X\ d)!i) = g\text{-}i\text{-}vec\ M\ X\ i\ (\text{degree } p + 1)$

<proof>

lemma *no-zeros-on-diagonal-helper:*

assumes *degree* $p \geq 1$

assumes $i < \text{degree } p$

assumes $M > 0$

assumes $X > 0$

shows $((\text{form-basis-helper } p\ M\ X\ (\text{degree } p - 1))!i)\$i = M * X^{\wedge} i$

<proof>

4.2 Casting lemmas

lemma *casted-distinct-is-distinct:*

fixes *vecs*:: *int vec list*

assumes *distinct-vecs*: *distinct vecs*

shows *distinct* $((\text{map of-int-vec } \text{vecs})::\text{rat vec list})$

<proof>

Copy pasted from VS_Connect, which only has it for field coefficients

lemma *(in vec-module) finsum-dim:*

shows $f \in A \rightarrow \text{carrier-vec } n \implies$

$\text{dim-vec } (\text{finsum } V\ f\ A) = n$

<proof>

lemma *(in vec-module) finsum-index:*

assumes $i: i < n$

and $f: f \in X \rightarrow \text{carrier-vec } n$

and $X: X \subseteq \text{carrier-vec } n$

shows $\text{finsum } V\ f\ X\ \$\ i = \text{sum } (\lambda x. f\ x\ \$\ i)\ X$

<proof>

lemma *is-int-rat-mul-of-int:*

assumes *snd* $(\text{quotient-of } y)\ \text{dvd } x$

shows *is-int-rat* $(\text{of-int } x * y)$

<proof>

lemma *casting-lin-comb-helper-set:*

assumes $vs: vs \subseteq \text{carrier-vec } n$
assumes $ldi: \text{module.lin-dep class-ring (module-vec TYPE(rat) } n)$
 $(\text{of-int-vec ' } vs)$
shows $\text{module.lin-dep class-ring (module-vec TYPE(int) } n) vs$
 $\langle \text{proof} \rangle$

lemma *casting-lin-comb-helper*:

assumes $\text{dim-vecs: } \bigwedge v. v \in \text{set vecs} \implies \text{dim-vec } v = \text{len-vec}$
assumes $\text{module.lin-indpt class-ring (module-vec TYPE(int) (len-vec)) (set vecs)}$
shows $\neg (\text{module.lin-dep class-ring (module-vec TYPE(rat) (len-vec))}$
 $(\text{set ((map of-int-vec vecs)::rat vec list))))$
 $\langle \text{proof} \rangle$

lemma *casting-lin-comb-helper-set-2*:

assumes $vs: vs \subseteq \text{carrier-vec } n$
assumes $ldi: \text{module.lin-dep class-ring (module-vec TYPE(int) } n) vs$
shows $\text{module.lin-dep class-ring (module-vec TYPE(rat) } n)$
 $(\text{of-int-vec ' } vs)$
 $\langle \text{proof} \rangle$

lemma *casting-lin-comb-helper-2*:

assumes $\text{dim-vecs: } \bigwedge v. v \in \text{set vecs} \implies \text{dim-vec } v = \text{len-vec}$
assumes $\neg (\text{module.lin-dep class-ring (module-vec TYPE(rat) (len-vec))}$
 $(\text{set ((map of-int-vec vecs)::rat vec list))))$
shows $\text{module.lin-indpt class-ring (module-vec TYPE(int) (len-vec)) (set vecs)}$
 $\langle \text{proof} \rangle$

4.3 Properties of associated polynomial

lemma *f-representation*:

shows $f = (\sum i < \text{degree } f. \text{monom (poly.coeff } f \ i) \ i) +$
 $\text{monom (lead-coeff } f) (\text{degree } f)$
 $\langle \text{proof} \rangle$

lemma *vec-associated-to-int-poly-inverse*:

assumes $X > 0$
fixes $f:: \text{int poly}$
shows $f = \text{int-poly-associated-to-vec (vec-associated-to-int-poly } f \ X) \ X$
 $\langle \text{proof} \rangle$

lemma *int-poly-associated-to-vec-degree-helper-le*:

shows $\text{degree } (\sum i \leq n. \text{monom (f } i) \ i) \leq n$
 $\langle \text{proof} \rangle$

lemma *int-poly-associated-to-vec-degree-helper-lt*:

assumes $n > 0$
shows $\text{degree } (\sum i < n. \text{monom (f } i) \ i) < n$
 $\langle \text{proof} \rangle$

lemma *int-poly-associated-to-vec-degree*:
fixes $v:: \text{int vec}$
assumes $\text{dim-vec } v > 0$
shows $\text{degree } (\text{int-poly-associated-to-vec } v \ X) < \text{dim-vec } v$
 $\langle \text{proof} \rangle$

lemma *degree-associated-poly*:
shows $\text{degree } (\text{int-poly-associated-to-vec } v \ X) \leq \text{dim-vec } v$
 $\langle \text{proof} \rangle$

lemma *degree-associated-poly-lt*:
assumes $X > 0$
assumes $\text{dim-vec } v \geq 1$
shows $\text{degree } (\text{int-poly-associated-to-vec } v \ X) < \text{dim-vec } v$
 $\langle \text{proof} \rangle$

lemma *degree-of-monom-sum-list*:
fixes $\text{ell}:: \text{int list}$
fixes $j:: \text{nat}$
assumes $\text{ell} \neq []$
shows degree
 $\quad (\sum_{i < \text{length } \text{ell}} \text{monom}$
 $\quad \quad (\text{ell } ! \ i) \ i) < \text{length } \text{ell}$
 $\langle \text{proof} \rangle$

lemma *coeff-of-monom-sum-list*:
fixes $\text{ell}:: \text{int list}$
fixes $j:: \text{nat}$
assumes $j < \text{length } \text{ell}$
shows coeff
 $\quad (\sum_{i < \text{length } \text{ell}} \text{monom}$
 $\quad \quad (\text{ell } ! \ i) \ i) \ j = \text{ell } ! \ j$
 $\langle \text{proof} \rangle$

lemma *coeff-of-monom-sum*:
fixes $a:: \text{int vec}$
assumes $i < \text{dim-vec } a$
shows coeff
 $\quad (\sum_{i < \text{dim-vec } a} \text{monom}$
 $\quad \quad (a \ \$ \ i) \ i) \ i = a \ \$ \ i$
 $\langle \text{proof} \rangle$

This next lemma required showing that every vector in the lattice has its i th component divisible by X^n (a property of the basis).

lemma *access-entry-in-int-poly-associated-to-vec*:
fixes $x \ i:: \text{nat}$
fixes $v:: \text{int vec}$
assumes $i < \text{dim-vec } v$
shows $(\text{coeff } (\text{int-poly-associated-to-vec } v \ X) \ i) = \lfloor \text{real-of-int } (v \ \$ \ i) / \text{real } (X$

$\hat{i}]$
 $\langle proof \rangle$

lemma *dim-vec-associated-to-int-poly-lteq*:

fixes $v:: int\ vec$
assumes $dim-vec\ v \geq 1$
assumes $X > 0$
shows $dim-vec\ (vec-associated-to-int-poly\ (int-poly-associated-to-vec\ v\ X)\ X) \leq dim-vec\ v$
 $\langle proof \rangle$

lemma *dim-vec-associated-to-int-poly-lt-imp-zeros*:

fixes $v:: int\ vec$
fixes $i:: nat$
assumes $entries-div: \bigwedge j. j < dim-vec\ v \implies (\exists k. (X \hat{j}) * k = (v \$ j))$
assumes $X-gt: X > 0$
assumes $dim-vec\ (vec-associated-to-int-poly\ (int-poly-associated-to-vec\ v\ X)\ X) < dim-vec\ v$
assumes $i-gt: i \geq dim-vec\ (vec-associated-to-int-poly\ (int-poly-associated-to-vec\ v\ X)\ X)$
assumes $i-lt: i < dim-vec\ v$
shows $v \$ i = 0$
 $\langle proof \rangle$

lemma *int-poly-associated-to-notNil-vec-same-norm-upto*:

assumes $dim-vec\ v \geq 1$
assumes $X-gt: X > 0$
assumes $entries-div: \bigwedge j. j < dim-vec\ v \implies (\exists k. (X \hat{j}) * k = (v \$ j))$
assumes $w = (vec-associated-to-int-poly\ (int-poly-associated-to-vec\ v\ X)\ X)$
shows $(\sum i < (dim-vec\ v). (v \$ i) \hat{2}) = (\sum i < (dim-vec\ w). (v \$ i) \hat{2})$
 $\langle proof \rangle$

lemma *int-poly-associated-to-notNil-vec-same-entries*:

fixes $i:: nat$
fixes $v\ w:: int\ vec$
assumes $i-lt: i < dim-vec\ w$
assumes $dim-vec\ v \geq 1$
assumes $X-gt: X > 0$
assumes $entries-div: \bigwedge j. j < dim-vec\ v \implies (\exists k. (X \hat{j}) * k = (v \$ j))$
assumes $w-is: w = (vec-associated-to-int-poly\ (int-poly-associated-to-vec\ v\ X)\ X)$
shows $((v \$ i) = (w \$ i))$
 $\langle proof \rangle$

lemma *int-poly-associated-to-nil-vec-same-norm*:

assumes $X > 0$
assumes $dim-vec\ v = 0$
shows $euclidean-norm-int-vec\ v = euclidean-norm-int-vec\ (vec-associated-to-int-poly\ (int-poly-associated-to-vec\ v\ X)\ X)$
 $\langle proof \rangle$

lemma *int-poly-associated-to-notNil-vec-same-norm*:
assumes $X\text{-gt}: X > 0$
assumes $\text{dim-vec } v > 0$
assumes $\bigwedge j. j < \text{dim-vec } v \implies (\exists k. (X \hat{=} j) * k = (v \$ j))$
shows $\text{euclidean-norm-int-vec } v = \text{euclidean-norm-int-vec } (\text{vec-associated-to-int-poly } (int\text{-poly-associated-to-vec } v X) X)$
 $\langle \text{proof} \rangle$

lemma *int-poly-associated-to-vec-same-norm*:
assumes $X > 0$
assumes $\bigwedge j. j < \text{dim-vec } v \implies (\exists k. (X \hat{=} j) * k = (v \$ j))$
shows $\text{euclidean-norm-int-vec } v = \text{euclidean-norm-int-vec } (\text{vec-associated-to-int-poly } (int\text{-poly-associated-to-vec } v X) X)$
 $\langle \text{proof} \rangle$

lemma *int-poly-associated-to-vec-sum*:
assumes $\text{dim-vec } a = \text{dim-vec } b$
assumes $\bigwedge i. i < \text{dim-vec } a \implies a \$ i \text{ mod } X \hat{=} i = 0$
assumes $\bigwedge i. i < \text{dim-vec } b \implies b \$ i \text{ mod } X \hat{=} i = 0$
assumes $X > 0$
shows $\text{int-poly-associated-to-vec } (a+b) X = \text{int-poly-associated-to-vec } a X + \text{int-poly-associated-to-vec } b X$
 $\langle \text{proof} \rangle$

lemma *int-poly-associated-to-vec-constant-mult*:
fixes $c :: \text{nat} \Rightarrow \text{int}$
fixes $v :: \text{int vec}$
assumes $X > 0$
assumes $\text{div-by-}X: \bigwedge i. i < \text{dim-vec } v \implies v \$ i \text{ mod } X \hat{=} i = 0$
shows $\text{int-poly-associated-to-vec } (c \text{ i } \cdot_v v) X = \text{smult } (c \text{ i}) (\text{int-poly-associated-to-vec } v X)$
 $\langle \text{proof} \rangle$

4.4 Generic Coppersmith assumptions locale

The generic properties required are stored in this locale.

locale *coppersmith-assms* = *LLL-with-assms* +
fixes $x0 :: \text{int}$
fixes $M X :: \text{nat}$
fixes $F :: \text{int poly}$
assumes $M: M > 0$
assumes $X: X > 0$
assumes $n: n > 0$
assumes $x0: \text{poly } F x0 \text{ mod } M = 0 \ |x0| \leq \text{int } X$
assumes $\text{lfs: length fs-init} \neq 0$

assumes $Xpoly: \bigwedge v j.$
 $v \in set\ fs-init \implies j < n \implies$
 $v \$ j\ mod\ int\ X \wedge j = 0$
assumes $rt: \bigwedge v.$
 $v \in set\ fs-init \implies$
 $poly\ (int-poly-associated-to-vec\ v\ X)\ x0\ mod\ M = 0$
begin

lemma *sumlist-mod*:
assumes $\bigwedge v. v \in set\ xs \implies dim-vec\ v = n$
assumes $m0: \bigwedge v j.$
 $v \in set\ xs \implies j < n \implies$
 $v \$ j\ mod\ int\ X \wedge j = 0$
assumes $i < n$
shows $M.sumlist\ xs\ \$\ i\ mod\ X \wedge i = 0$
 $\langle proof \rangle$

lemma *poly-associated-to-vec-sumlist*:
assumes $\bigwedge v. v \in set\ xs \implies dim-vec\ v = n$
assumes $\bigwedge v j. v \in set\ xs \implies$
 $j < dim-vec\ v \implies v \$ j\ mod\ X \wedge j = 0$
shows $int-poly-associated-to-vec\ (sumlist\ xs)\ X =$
 $(\sum\ i < length\ xs. int-poly-associated-to-vec\ (xs\ !\ i)\ X)$
 $\langle proof \rangle$

lemma *short-vector-inherit-props*:
shows $\bigwedge j. j < n \implies short-vector\ \$\ j\ mod\ X \wedge j = 0$
 $poly\ (int-poly-associated-to-vec\ short-vector\ X)\ x0\ mod\ M = 0$
 $\langle proof \rangle$

lemma *root-poly-short-vector*:
assumes $bnd: real-of-int\ \|short-vector\|^2 < M^2 / n$
shows $poly\ (int-poly-associated-to-vec\ short-vector\ X)\ x0 = 0$
 $\langle proof \rangle$

lemma *bnd-raw-imp-short-vec-bound*:
assumes $bnd-raw:$
 $(real-of-rat\ \alpha) \wedge (m * (m - 1) \div 2) *$
 $real-of-int\ (gs.Gramian-determinant\ fs-init\ m) *$
 $(real\ n) \wedge m <$
 $M \wedge (2 * m)$
shows $real-of-int\ \|short-vector\|^2 < M^2 / n$
 $\langle proof \rangle$

end

end

5 Proof of Lightweight Algorithm

We start by proving the "lightweight algorithm" as a stepping stone to the full algorithm (proved in `Coppersmith.thy`).

theory *Towards-Coppersmith*

imports *Coppersmith-Generic*
begin

5.1 Basic properties

lemma *dim-form-basis-helper*:

shows $\text{length } (\text{form-basis-helper } p \ M \ X \ d) = d + 1$
<proof>

lemma *dim-form-basis*:

shows $\text{length } (\text{form-basis } p \ M \ X \ d) = d + 2$
<proof>

5.2 Matrix properties

5.2.1 Basic properties and preliminaries

lemma *dim-row-basis-mat*:

assumes $\text{degree } p \geq 1$
shows $\text{dim-row } (\text{vec-list-to-square-mat } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1))) = \text{degree } p + 1$
<proof>

lemma *dim-col-basis-mat*:

assumes $\text{degree } p \geq 1$
shows $\text{dim-col } (\text{vec-list-to-square-mat } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1))) = (\text{degree } p + 1)$
<proof>

lemma *matrix-carrier*:

assumes $\text{degree } p \geq 1$
assumes $i < \text{degree } p + 1$
shows $\text{vec-list-to-square-mat } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1)) \in \text{carrier-mat } (\text{degree } p + 1) \ (\text{degree } p + 1)$
<proof>

lemma *vector-sum-monom*:

fixes $v:: \text{int } \text{vec}$
fixes $d \ i:: \text{nat}$
assumes $d = \text{dim-vec } v$
assumes $d > 0$
assumes $i\text{-lt}: i < d$
assumes $\text{zero-monom}: \bigwedge j::\text{nat}. j < d \implies j \neq i \implies \text{monom } (v \ \$ \ j) \ j = 0$

shows $(\sum_{i < d} \text{monom } (v \$ i) i) = \text{monom } (v \$ i) i$
 <proof>

lemma *int-poly-associated-to-g-i-vec*:

assumes *X-gt*: $X > 0$

assumes *M-gt*: $M > 0$

assumes *i-lt*: $i \leq (\text{degree } p)$

shows *int-poly-associated-to-vec* $(g\text{-i-vec } M X i (\text{degree } p + 1)) X = \text{monom } M i$
 <proof>

lemma *ith-row-form-basis*:

shows $i \leq d \implies ((\text{form-basis } p M X d)!i) = (\text{form-basis-helper } p M X d)!i$

$((\text{form-basis } p M X d)!(d+1)) = \text{vec-associated-to-int-poly } p X$

<proof>

lemma *set-form-basis*:

shows $x \in \text{set } (\text{form-basis } p M X (\text{degree } p - 1)) \implies x = \text{vec-associated-to-int-poly } p X \vee$

$x \in \text{set } (\text{form-basis-helper } p M X (\text{degree } p - 1))$

<proof>

lemma *dim-vector-in-basis*:

fixes *i*:: *nat*

assumes $i < d + 2$

shows *dim-vec* $((\text{form-basis } p M X d)! i) = (\text{degree } p+1)$

<proof>

5.2.2 Properties of matrix associated to input

lemma *matrix-row-form-basis-carrier*:

assumes *degree* $p \geq 1$

assumes $i < \text{degree } p + 1$

shows *form-basis* $p M X (\text{degree } p - 1)! i \in \text{carrier-vec } (\text{degree } p + 1)$

<proof>

lemma *matrix-row-form-basis*:

assumes *degree* $p \geq 1$

assumes *i-lt*: $i < \text{degree } p + 1$

shows *row* $(\text{vec-list-to-square-mat } (\text{form-basis } p M X (\text{degree } p - 1))) i =$
 $(\text{form-basis } p M X (\text{degree } p - 1)! i)$

<proof>

lemma *matrix-diagonal-element*:

assumes *degree* $p \geq 1$

assumes $i < \text{degree } p$

shows *vec-list-to-square-mat*

$(\text{form-basis } p M X (\text{degree } p - 1)) \$\$$

$(i, i) = ((\text{form-basis } p M X (\text{degree } p - 1)! i)\i

<proof>

lemma *no-zeros-on-diagonal*:

assumes $\text{degree } p \geq 1$
assumes $M > 0$
assumes $X > 0$
shows $0 \notin \text{set } (\text{diag-mat } (\text{vec-list-to-square-mat } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1))))$
<proof>

lemma *form-basis-helper-is-lower-triangular*:

fixes $i \ j :: \text{nat}$
assumes $i < j$
assumes $j < (\text{degree } p + 1)$
shows $((\text{form-basis-helper } p \ M \ X \ (\text{degree } p - 1))!i)!j = 0$
<proof>

lemma *form-basis-is-lower-triangular*:

fixes $i \ j :: \text{nat}$
assumes $i\text{-lt}: i < j$
assumes $j\text{-lt}: j < (\text{degree } p + 1)$
shows $(\text{vec-list-to-square-mat } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1))) \ \$\$ \ (i,j) = 0$
<proof>

lemma *form-basis-distinct*:

assumes $\text{degree } p \geq 1$
assumes $M > 0$
assumes $X > 0$
shows $\text{distinct } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1))$
<proof>

lemma *det-of-matrix*:

fixes $M \ X :: \text{nat}$
assumes $M > 0$
assumes $X > 0$
assumes $\text{degree } p \geq 1$
assumes $d\text{-is}: d = \text{degree } p$
assumes $n\text{-is}: n = (\sum i \leq d. i)$
assumes $\text{monic-poly}: \text{coeff } p \ d = 1$
shows $\text{det } (\text{vec-list-to-square-mat } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1))) = M^{\wedge}(\text{degree } p) * X^{\wedge}n$
<proof>

5.2.3 Properties of casted matrix

lemma *dim-row-basis-of-int-mat*:

assumes $\text{degree } p \geq 1$
shows $\text{dim-row } (\text{vec-list-to-square-mat } (\text{map of-int-vec } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1)))) = \text{degree } p + 1$
 ⟨proof⟩

lemma *dim-col-basis-of-int-mat:*

assumes $\text{degree } p \geq 1$
shows $\text{dim-col } (\text{vec-list-to-square-mat } (\text{map of-int-vec } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1)))) = (\text{degree } p + 1)$
 ⟨proof⟩

lemma *of-int-matrix-row-form-basis-carrier:*

assumes $\text{degree } p \geq 1$
assumes $i < \text{degree } p + 1$
shows $(\text{map of-int-vec } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1))) ! i \in \text{carrier-vec } (\text{degree } p + 1)$
 ⟨proof⟩

lemma *of-int-matrix-carrier:*

assumes $\text{degree } p \geq 1$
assumes $i < \text{degree } p + 1$
shows $\text{vec-list-to-square-mat } (\text{map of-int-vec } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1))) \in \text{carrier-mat } (\text{degree } p + 1) \ (\text{degree } p + 1)$
 ⟨proof⟩

lemma *of-int-matrix-row-form-basis:*

assumes $\text{degree } p \geq 1$
assumes $i\text{-lt: } i < \text{degree } p + 1$
shows $\text{row } (\text{vec-list-to-square-mat } (\text{map of-int-vec } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1)))) i = (\text{map of-int-vec } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1))) ! i$
 ⟨proof⟩

lemma *of-int-matrix-row-form-basis-var:*

assumes $\text{degree } p \geq 1$
assumes $i\text{-lt: } i < \text{degree } p + 1$
shows $\text{row } (\text{vec-list-to-square-mat } (\text{map of-int-vec } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1)))) i = \text{of-int-vec } ((\text{form-basis } p \ M \ X \ (\text{degree } p - 1)) ! i)$
 ⟨proof⟩

lemma *of-int-mat-element:*

fixes $i \ j :: \text{nat}$
assumes $\text{degree } p \geq 1$
assumes $i < (\text{degree } p + 1)$
assumes $j < (\text{degree } p + 1)$
shows $(\text{vec-list-to-square-mat } (\text{map of-int-vec } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1)))) \$\$ (i,j) = \text{of-int } ((\text{vec-list-to-square-mat } (\text{form-basis } p \ M \ X \ (\text{degree } p - 1))) \$\$ (i,j))$
 ⟨proof⟩

lemma *of-int-mat-is-lower-triangular:*

fixes $i\ j::\text{nat}$
assumes $\text{degree } p \geq 1$
assumes $i < j$
assumes $j < (\text{degree } p + 1)$
shows $(\text{vec-list-to-square-mat } (\text{map of-int-vec } (\text{form-basis } p\ M\ X\ (\text{degree } p - 1))))\ \$(i,j) = 0$
<proof>

lemma *of-int-mat-no-zeros-on-diagonal:*

assumes $p\text{-gt: } \text{degree } p \geq 1$
assumes $M > 0$
assumes $X > 0$
shows $0 \notin \text{set } (\text{diag-mat } (\text{vec-list-to-square-mat } ((\text{map of-int-vec } (\text{form-basis } p\ M\ X\ (\text{degree } p - 1))))::\text{rat vec list}))$
<proof>

lemma *of-int-mat-form-basis-distinct:*

assumes $\text{degree } p \geq 1$
assumes $M > 0$
assumes $X > 0$
shows $\text{distinct } ((\text{map of-int-vec } (\text{form-basis } p\ M\ X\ (\text{degree } p - 1))))::\text{rat vec list}$

<proof>

lemma *of-int-form-basis-lin-ind:*

assumes $M > 0$
assumes $X > 0$
assumes $\text{degree } p \geq 1$
shows $\neg \text{module.lin-dep class-ring } (\text{module-vec } \text{TYPE}(\text{rat})\ (\text{degree } p + 1))$
 $(\text{set } ((\text{map of-int-vec } (\text{form-basis } p\ M\ X\ (\text{degree } p - 1))))::\text{rat vec list})$
<proof>

5.3 Top-level proof

lemma *towards-coppersmith:*

fixes $p\ f::\ \text{int poly}$
fixes $M\ X::\ \text{nat}$
fixes $x0::\ \text{int}$
assumes $\text{zero-mod-M: } \text{poly } p\ x0\ \text{mod } M = 0$
assumes $d\text{-is: } d = \text{degree } p$
assumes $d\text{-gt: } d > 0$
assumes $\text{monic-poly: } \text{coeff } p\ d = 1$
assumes $X\text{-gt: } X > 0$
assumes $X\text{-lt: } X < 1/(\text{sqrt } 2)*1/(\text{root } d\ (d+1))*(\text{root } (d*(d+1))\ M) \sim 2$
assumes $M\text{-gt: } M > 0$
assumes $x0\text{-le: } \text{abs } x0 \leq X$
assumes $f\text{-is: } f = \text{towards-coppersmith } p\ M\ X$
shows $\text{poly } f\ x0 = 0$

<proof>

lemma *towards-coppersmith-pretty*:

fixes $p f:: int\ poly$

fixes $M X:: nat$

fixes $x0:: int$

defines $d \equiv degree\ p$

defines $f \equiv towards-coppersmith\ p\ M\ X$

assumes *monic-poly*: $monic\ p$

assumes $d > 0$ **and** $M > 0$ **and** $X > 0$

assumes *zero-mod-M*: $poly\ p\ x0\ mod\ M = 0$

assumes *X-lt*: $X < 1/(sqrt\ 2)*1/(root\ d\ (d+1))*(root\ (d*(d+1))\ M)^2$

assumes *x0-le*: $abs\ x0 \leq X$

shows $poly\ f\ x0 = 0$

<proof>

end

6 Proof of Coppersmith's Method

In this file, we prove the full version of Coppersmith's method.

theory *Coppersmith*

imports *Coppersmith-Generic*

begin

6.1 Preliminaries and setup

lemma *calculate-h-coppersmith-aux-gteq1*:

fixes $e::real$

assumes $degree\ p > 1$

assumes $e > 0$

shows $calculate-h-coppersmith-aux\ p\ e \geq 1$

<proof>

lemma *calculate-h-coppersmith-aux-gt1*:

assumes *deg-gt*: $degree\ p > 1$

assumes $e > 0$

assumes *e-lt2*: $e < 1/(real\ (degree\ p))$

shows $calculate-h-coppersmith-aux\ p\ e > 1$

<proof>

6.1.1 Dimension properties of matrix

lemma *dim-vector-vec-associated-to-int-poly-padded*:

shows $dim-vec\ (vec-associated-to-int-poly-padded\ n\ p\ X) = n$

<proof>

lemma *dim-vector-row-of-coppersmith-matrix*:

shows $dim-vec\ (row-of-coppersmith-matrix\ p\ M\ X\ h\ i\ j) = (degree\ p)*h$

<proof>

lemma *dim-vector-form-basis-coppersmith-aux:*

fixes *i:: nat*

assumes *i < (degree p)*

shows *dim-vec ((form-basis-coppersmith-aux p M X h j) ! i) = (degree p)*h*

<proof>

lemma *length-form-basis-coppersmith-aux:*

shows *length (form-basis-coppersmith-aux p M X h j) = degree p*

<proof>

lemma *length-form-basis-coppersmith:*

fixes *h:: nat*

assumes *h-gt: h > 0*

shows *length (form-basis-coppersmith p M X h) = (degree p)*h*

<proof>

lemma *concat-property-helper:*

assumes *j < h*

shows *concat (map ($\lambda i. f i$) [0.. h]) = concat (map ($\lambda i. f i$) [0.. j]) @ concat (map ($\lambda i. f i$) [j .. h])*

<proof>

lemma *concat-equal-lists-length:*

fixes *f:: nat \Rightarrow int vec list*

fixes *i j d:: nat*

assumes *len-is: $\bigwedge i. i < h \implies \text{length } (f i) = d$*

shows *length (concat (map ($\lambda i. f i$) [0.. h])) = $d*h$*

<proof>

lemma *concat-property:*

fixes *f:: nat \Rightarrow int vec list*

fixes *i j d:: nat*

assumes *h > 0*

assumes *d-gt: d > 0*

assumes *r-lt: r < $d*h$*

assumes *len-is: $\bigwedge i. i < h \implies \text{length } (f i) = d$*

assumes *j-eq: $j = \text{nat } \lfloor \text{real } r / \text{real } d \rfloor$*

assumes *i-eq: $i = r - d * j$*

shows *(concat (map ($\lambda i. f i$) [0.. h])) ! r = (f j) ! i*

<proof>

lemma *row-of:*

assumes *r-lt: r < (degree p)*h*

defines *d \equiv degree p*

defines *j \equiv nat $\lfloor \text{real } r / \text{real } d \rfloor$*

defines *i \equiv r - d * j*

shows *(form-basis-coppersmith p M X h) ! r =*

$((\text{form-basis-coppersmith-aux } p \ M \ X \ h \ j) ! i)$
 $\langle \text{proof} \rangle$

lemma *dim-vector-form-basis-coppersmith*:

fixes $i :: \text{nat}$
assumes $i < (\text{degree } p) * h$
shows $\text{dim-vec } ((\text{form-basis-coppersmith } p \ M \ X \ h) ! i) = (\text{degree } p) * h$
 $\langle \text{proof} \rangle$

6.1.2 Equivalent Coppersmith matrix

definition *form-coppersmith-matrix*:: $\text{int poly} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{int mat}$

where $\text{form-coppersmith-matrix } p \ M \ X \ h = \text{mat } ((\text{degree } p) * h) ((\text{degree } p) * h)$
 $(\lambda(r, c). (\text{let } d = \text{degree } p; j = \text{nat } (\text{floor } (r/d)); i = (r - d*j) \text{ in } (M^{(h-1-j)}) * (\text{coeff } (p^{j*} (\text{monom } 1 \ i)) \ c) * X^c))$

lemma *matrix-match*:

assumes $r\text{-lt}: r < (\text{degree } p) * h$
assumes $c\text{-lt}: c < (\text{degree } p) * h$
assumes $h > 0$
shows $(\text{vec-list-to-square-mat } (\text{form-basis-coppersmith } p \ M \ X \ h)) \$\$ (r, c) = (\text{form-coppersmith-matrix } p \ M \ X \ h) \$\$ (r, c)$
 $\langle \text{proof} \rangle$

6.1.3 Lower triangular

lemma *form-coppersmith-matrix-is-lower-triangular*:

fixes $r \ c :: \text{nat}$
assumes $h > 0$
assumes $r\text{-lt}: r < c$
assumes $c\text{-lt}: c < (\text{degree } p) * h$
shows $(\text{form-coppersmith-matrix } p \ M \ X \ h) \$\$ (r, c) = 0$
 $\langle \text{proof} \rangle$

lemma *form-coppersmith-basis-is-lower-triangular*:

fixes $i \ j :: \text{nat}$
assumes $h > 0$
assumes $i\text{-lt}: i < j$
assumes $j\text{-lt}: j < (\text{degree } p) * h$
shows $(\text{vec-list-to-square-mat } (\text{form-basis-coppersmith } p \ M \ X \ h)) \$\$ (i, j) = 0$
 $\langle \text{proof} \rangle$

6.1.4 Distinct elements

lemma *coppersmith-matrix-carrier-mat*:

assumes $h > 0$
shows $\text{vec-list-to-square-mat } (\text{form-basis-coppersmith } p \ M \ X \ h) \in \text{carrier-mat } ((\text{degree } p) * h) ((\text{degree } p) * h)$
 $\langle \text{proof} \rangle$

lemma *no-zeros-on-diagonal-coppersmith*:

assumes *degree* $p \geq 1$
assumes *M-gt*: $M > 0$
assumes *X-gt*: $X > 0$
assumes *h-gt*: $h > 0$
shows $0 \notin \text{set} (\text{diag-mat} (\text{vec-list-to-square-mat} (\text{form-basis-coppersmith } p \ M \ X \ h)))$
<proof>

lemma *form-basis-coppersmith-distinct*:

fixes $M \ X :: \text{nat}$
assumes $1 \leq \text{degree } p$
assumes *p-neg*: $p \neq 0$
assumes *M-gt*: $M > 0$
assumes *X-gt*: $X > 0$
assumes *h-gt*: $h > 0$
shows *distinct* (*form-basis-coppersmith* $p \ M \ X \ h$)
<proof>

lemma *matrix-row-form-basis-coppersmith*:

assumes *degree* $p \geq 1$
assumes *i-lt*: $i < (\text{degree } p) * h$
assumes $h > 0$
shows *row* (*vec-list-to-square-mat* (*form-basis-coppersmith* $p \ M \ X \ h$)) $i = (\text{form-basis-coppersmith } p \ M \ X \ h) ! i$
<proof>

6.1.5 Linear independence properties

lemma *form-basis-coppersmith-lin-ind*:

assumes $M > 0$
assumes $X > 0$
assumes *degree* $p \geq 1$
assumes $h > 0$
shows $\neg \text{module.lin-dep class-ring} (\text{module-vec } \text{TYPE}(int) ((\text{degree } p) * h))$
 $(\text{set} (\text{form-basis-coppersmith } p \ M \ X \ h))$
<proof>

6.1.6 Divisible by X property

lemma *row-of-cs-matrix-div-by-Xpow*:

fixes $M \ X \ i \ j \ h :: \text{nat}$
assumes *i-lt*: $i < \text{degree } p$
assumes *j-lt*: $j < h$
assumes *j-lt*: $ja < h * (\text{degree } p)$
shows (*row-of-coppersmith-matrix* $p \ M \ X \ h \ i \ j$) $\$ ja \text{ mod } (X \wedge ja) = 0$
<proof>

lemma *form-basis-coppersmith-div-by-Xpow*:

fixes $M \ X :: \text{nat}$

fixes a : *int vec*
assumes j -lt: $ja < h * (\text{degree } p)$
assumes a -inset: $a \in \text{set } (\text{form-basis-coppersmith } p \ M \ X \ h)$
shows $(a \ \$ \ ja) \ \text{mod } (X \wedge ja) = 0$
 <proof>

6.1.7 Zero mod M property

lemma *row-of-cs-matrix-zero-mod-M*:

fixes $M \ X \ i \ j \ h$::*nat*
assumes p -neq: $p \neq 0$
assumes M -gt: $M > 0$
assumes X -gt: $X > 0$
assumes h -gt1: $h > 1$
assumes p -zero-mod-M: $\text{poly } p \ x0 \ \text{mod } M = 0$
assumes deg-gt : $\text{degree } p > 1$
assumes i -lt: $i < \text{degree } p$
assumes j -lt: $j < h$
shows $\text{poly } (\text{int-poly-associated-to-vec } (\text{row-of-coppersmith-matrix } p \ M \ X \ h \ i \ j) \ X) \ x0 \ \text{mod } M \wedge^{(h-1)} = 0$
 <proof>

lemma *form-basis-coppersmith-zero-mod-M*:

fixes $M \ X$::*nat*
assumes p -neq: $p \neq 0$
assumes M -gt: $M > 0$
assumes X -gt: $X > 0$
assumes h -gt: $h > 1$
assumes p -zero-mod-M: $\text{poly } p \ x0 \ \text{mod } M = 0$
assumes a -inset: $a \in \text{set } (\text{form-basis-coppersmith } p \ M \ X \ h)$
assumes deg-gt : $\text{degree } p > 1$
shows $\text{poly } (\text{int-poly-associated-to-vec } a \ X) \ x0 \ \text{mod } M \wedge^{(h-1)} = 0$
 <proof>

6.1.8 Determinant of matrix

lemma *determinant-bound-arithmetic-helper*:

fixes k ::*nat*
shows $(\prod_{j < (w+1)}. k \wedge^j) = \text{sqrt } (k \wedge^{(w * (w+1))})$
 <proof>

lemma *det-of-form-coppersmith-matrix*:

fixes $M \ X$::*nat*
assumes $M > 0$
assumes $X > 0$
assumes d -is: $d = \text{degree } p$
assumes monic-poly : $\text{coeff } p \ d = 1$
assumes h -gt: $h > 1$
assumes deg-gt : $\text{degree } p > 1$
shows $\text{det } (\text{form-coppersmith-matrix } p \ M \ X \ h) =$

($\text{root } 2 (M^{(h-1)*d*h})$) * ($\text{root } 2 (X^{(d*h-1)*d*h})$)
 <proof>

lemma *det-of-matrix*:

fixes $M X:: \text{nat}$
assumes $M > 0$
assumes $X > 0$
assumes $d > 1$
assumes $d\text{-is}$: $d = \text{degree } p$
assumes monic-poly : $\text{coeff } p \ d = 1$
assumes $h\text{-gt}$: $h > 1$
shows $\text{det } (\text{vec-list-to-square-mat } (\text{form-basis-coppersmith } p \ M \ X \ h)) =$
 $(\text{root } 2 (M^{(h-1)*d*h})) * (\text{root } 2 (X^{(d*h-1)*d*h}))$
 <proof>

6.2 Top-level proof

definition *root-bound*:: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{real} \Rightarrow \text{real}$
where $\text{root-bound } M \ d \ \text{eps} = 1/2*(M \ \text{powr } (1/d - \text{eps}))$

6.2.1 Arithmetic

lemma *epsilon-bounded-below*:

assumes $d > 0$
assumes $\text{eps} > 0$
assumes $d*h-1 > 0$
assumes $d = \text{degree } p$
assumes $h = \text{calculate-h-coppersmith } p \ \text{eps}$
shows $\text{eps} \geq (d-1)/(d*(d*h-1))$
 <proof>

lemma *z-arith*:

assumes $x: (x::\text{real}) \geq 0$
shows $x / (1 + x) \leq \ln (1 + x)$
 <proof>

lemma *powr-divide-both*:

assumes $(a::\text{real}) \geq 0 \ x > 0 \ b \ \text{powr } y \geq 1$
assumes le : $a \ \text{powr } x \geq b \ \text{powr } y$
shows $a \geq b \ \text{powr } (y / x)$
 <proof>

lemma *coppersmith-arithmetic-convergence-1*:

fixes $y:: \text{real}$
assumes $y: y \geq 1 / 0.18$
shows $0 \leq 1.414 - (1 + y) \ \text{powr } (1 / y)$
 <proof>

lemma *coppersmith-arithmetic-convergence*:

fixes $x:: \text{real}$

```

assumes  $x: 0 < x \leq 0.18$ 
shows  $(1 + (1/x)) \text{ powr } (x) \leq \text{sqrt } 2$ 
<proof>

```

6.2.2 Main results

```

lemma coppersmith-finds-small-roots:
fixes  $p f:: \text{int poly}$ 
fixes  $M X:: \text{nat}$ 
fixes  $x0:: \text{int}$ 
fixes  $\text{eps}:: \text{real}$ 
assumes zero-mod-M:  $\text{poly } p \text{ } x0 \text{ mod } M = 0$ 
assumes d-is:  $d = \text{degree } p$ 
assumes d-gt:  $d > 1$ 
assumes monic-poly:  $\text{coeff } p \text{ } d = 1$ 
assumes X-lt:  $X < \text{root-bound } M \text{ } d \text{ } \text{eps}$ 
assumes M-gt:  $M > 0$ 
assumes X-gt-zero:  $X > 0$ 
assumes x0-le:  $\text{abs } x0 \leq X$ 
assumes eps-le:  $\text{eps} \leq 0.18 * (d-1) / d$ 
assumes eps-lt:  $\text{eps} < 1 / (\text{real } (\text{degree } p))$ 
assumes eps-gt:  $\text{eps} > 0$ 
assumes f-is:  $f = \text{coppersmith } p \text{ } M \text{ } X \text{ } \text{eps}$ 
shows  $\text{poly } f \text{ } x0 = 0$ 
<proof>

```

```

theorem coppersmith-finds-small-roots-pretty:
fixes  $p f:: \text{int poly}$ 
fixes  $M X:: \text{nat}$ 
fixes  $x0:: \text{int}$ 
fixes  $\text{eps}:: \text{real}$ 
defines  $d \equiv \text{degree } p$ 
defines  $f \equiv \text{coppersmith } p \text{ } M \text{ } X \text{ } \text{eps}$ 
assumes monic-poly: monic  $p$ 
assumes  $d > 1$  and  $M > 0$  and  $X > 0$ 
assumes zero-mod-M:  $\text{poly } p \text{ } x0 \text{ mod } M = 0$ 
assumes X-lt:  $X < \text{root-bound } M \text{ } d \text{ } \text{eps}$ 
assumes x0-le:  $\text{abs } x0 \leq X$ 
assumes eps-le:  $\text{eps} \leq 0.18 * (d-1) / d$ 
assumes eps-lt:  $\text{eps} < 1 / d$ 
assumes eps-gt0:  $\text{eps} > 0$ 
shows  $\text{poly } f \text{ } x0 = 0$ 
<proof>

```

end

7 Examples of Coppersmith's Method

In this file, we provide some examples of Coppersmith's method, with correctness proofs (when applicable).

```
theory Coppersmith-Examples
```

```
imports Coppersmith  
        Towards-Coppersmith  
begin
```

7.1 Example of lightweight method

Following Example 19.1.6 in Galbraith. This example produces $[-444, 1, -20, -2]$, which corresponds to $-2x^3 - 20x^2 + x + 444$, which has 4 as a root. Computing $4^3 + 10 * 4^2 + 5000 * 4 - 222$ produces 20002, which is 0 mod 10001. Note that here, we cannot use our top-level correctness result for the lightweight method to prove correctness. This is because the conditions of this top-level result are not satisfied; however, the method succeeds despite not fulfilling the conditions of this result, because the LLL algorithm can be better than the bound in the result. See Example 19.1.6 in "Mathematics of Public Key Cryptography" by Galbraith for more discussion of this.

```
value towards-coppersmith  $[-222, 5000, 10, 1:]$  10001 10
```

7.2 Examples of Coppersmith's method

Following Exercise 19.1.12 in Galbraith.

This next example produces $15955575444164700778296 - 86948462676890416832x + 50262448764961319x^2 + 334700479564525x^3 - 611446097378x^4 - 577363178x^5 + 1008850x^6 + 8592x^7$, which has 267 as a root, which is a root of the original polynomial mod 2^9 .

```
value coppersmith  $[:227, 46976195, 2^{25}-2883584, 1:]$   $((2^{20} + 7)*(2^{21} + 17))$   $(2^9)$  0.089
```

We now prove that this example satisfies the conditions of our top-level correctness theorem for Coppersmith's method.

```
lemma coppersmith-finds-small-roots-example1:
```

```
  fixes  $p f:: int poly$   
  fixes  $M X:: nat$   
  fixes  $x0:: int$   
  fixes  $k:: nat$   
  defines  $p \equiv [:227, 46976195, 2^{25}-2883584, 1:]$   
  defines  $d \equiv degree\ p$   
  defines  $M \equiv ((2^{20} + 7)*(2^{21} + 17))$   
  defines  $X \equiv 2^9$   
  defines  $f \equiv coppersmith\ p\ M\ X\ 0.089$   
  assumes  $x0-le: |x0| \leq X$ 
```

```

assumes zero-mod-M: poly p x0 mod M = 0
shows poly f x0 = 0
⟨proof⟩

```

In this next example, we are trying to find small roots less than 3 of $x^3 + 1000x^2 + 25 + 1 \pmod{4059}$. Running "coppersmith" on this example yields `[41858, - 28457, 6100, 376, 150, - 31, - 91, - 28, - 17:]`, which corresponds to $-17x^8 - 28x^7 - 91x^6 - 31x^5 + 150x^4 + 376x^3 + 6100x^2 - 28457x + 41858$, which has 2 as a root. Plugging in $2^3 + 1000 * 2^2 + 25 * 2 + 1$ indeed yields 4059, which is 0 mod 4059.

```

value form-basis-coppersmith [1, 25, 1000, 1:] 4059 3 (calculate-h-coppersmith
[1, 25, 1000, 1:] 0.10)
value reduce-basis 2 (form-basis-coppersmith [1, 25, 1000, 1:] 4059 3 (calculate-h-coppersmith
[1, 25, 1000, 1:] 0.10))
value coppersmith [1, 25, 1000, 1:] 4059 3 0.10

```

We now prove that this example satisfies the conditions of our top-level correctness theorem for Coppersmith's method.

```

lemma coppersmith-finds-small-roots-example2:
  fixes p f:: int poly
  fixes M X:: nat
  fixes x0:: int
  fixes k:: nat
  defines p ≡ [1, 25, 1000, 1:]
  defines d ≡ degree p
  defines M ≡ 4059
  defines X ≡ 3
  defines f ≡ coppersmith p M X 0.10
  assumes x0-le: |x0| ≤ X
  assumes zero-mod-M: poly p x0 mod M = 0
  shows poly f x0 = 0
⟨proof⟩

```

end

References

- [1] J. Alperin-Sheriff and C. Peikert. Lattices in cryptography: Lecture 4, Coppersmith, cryptanalysis. Georgia Tech lecture notes, Fall 2023. Online lecture notes available at <https://web.eecs.umich.edu/~cpeikert/lic13/lec04.pdf>.
- [2] R. Bottesch, J. Divasón, M. W. Haslbeck, S. J. C. Joosten, R. Thiemann, and A. Yamada. A verified LLL algorithm. *Archive of Formal Proofs*, February 2018. https://isa-afp.org/entries/LLL_Basis_Reduction.html, Formal proof development.

- [3] S. D. Galbraith. *Mathematics of Public Key Cryptography*. Cambridge University Press, 2012.
- [4] W. Trappe and L. C. Washington. *Introduction to Cryptography with Coding Theory (2nd Edition)*. Prentice-Hall, Inc., USA, 2005.