

Conway's Circle Theorem in Isabelle/HOL

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Abstract

We formalize Conway's circle theorem in Isabelle/HOL. For a non-degenerate Euclidean triangle ABC with side lengths $a = |BC|$, $b = |CA|$, and $c = |AB|$, extend the two side-lines meeting at each vertex past that vertex by the length of the side opposite that vertex: at A both adjacent sides are extended by a , at B by b , and at C by c . The six resulting endpoints lie on one circle. Its center is the incenter I of the triangle, and its radius is

$$\sqrt{r^2 + s^2},$$

where r is the inradius and $s = (a + b + c)/2$ is the semiperimeter.

We give the result in three equivalent forms: as six explicit distance equations, as a subset relation $\subseteq$ into the corresponding **sphere**, and as a single bounded-universal statement over the six-point set. A separate construction-correctness theorem records that each Conway point lies on the intended side line and at the declared extension length from the adjacent vertex.

AI assistance was used for proof engineering.

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1 Introduction

Conway's construction is a compact Euclidean configuration: a triangle determines three side lengths, those lengths determine six points on the three side lines, and the resulting endpoints unexpectedly form a circle. The theorem is often presented as a visual fact about equal segments; the formal proof has to make explicit which directed lines are used, which extensions are external to the triangle, and how the metric computation is organized.

The development follows the elementary proof. Each constructed point lies on one of the three side lines. The perpendicular from the incenter I to that side has length r , and the signed distance from the foot of the perpendicular to either constructed endpoint is s . The Pythagorean theorem therefore gives the common distance

$$\sqrt{r^2 + s^2}$$

from I to all six endpoints.

2 Definitions

The formal statement works in the Euclidean plane over real coordinates. The vertices A , B , and C are assumed not to be collinear. The side lengths are named in the usual way:

$$\begin{aligned} a &= |BC|, & b &= |CA|, & c &= |AB|, \\ s &= \frac{a + b + c}{2}. \end{aligned}$$

The incenter I is the point whose perpendicular distance to each side line is the common inradius r .

For the construction, each side line contributes two points. On the line AB , the point A_c is beyond A with $|AA_c| = a$, and the point B_c is beyond B with $|BB_c| = b$. Cyclically, the line AC gives A_b and C_b , and the line BC gives B_a and C_a . These prose names are only a guide to the geometric objects; the precise predicates and uniqueness facts are those rendered from the Isabelle theory below.

3 Proof Sketch

Let T_{AB} be the touch point of the incircle with the side AB . The standard tangent-length identities give

$$|AT_{AB}| = s - a, \quad |BT_{AB}| = s - b.$$

Since A_c lies past A by length a , the directed distance from T_{AB} to A_c has absolute value s ; similarly the directed distance from T_{AB} to B_c has absolute

value s . The segment IT_{AB} is perpendicular to AB and has length r , so the two right triangles with hypotenuses IA_c and IB_c both have legs r and s . Thus

$$|IA_c| = |IB_c| = \sqrt{r^2 + s^2}.$$

The same argument on the side lines AC and BC proves the identical distance formula for A_b , C_b , B_a , and C_a . Hence all six constructed points lie on the circle with center I and radius $\sqrt{r^2 + s^2}$.

4 Historical Remark

The result is part of the informal triangle-geometry tradition associated with John H. Conway, and public expositions record the construction under his name [3]. We do not attempt a precise priority claim here. Later conic work by Evans fits the theorem into the broader triangle-center literature [2], and García Capitán subsequently showed that the same six-point pattern belongs to a family of conics, with the Conway circle recovered when the chosen point is the incenter [1]. Honsberger's collection gives related Euclidean-geometry background [4].

theory *Conway-Circle*

imports *HOL-Analysis.Analysis*

begin

definition *triangle-semiperimeter* :: $real \wedge 2 \Rightarrow real \wedge 2 \Rightarrow real \wedge 2 \Rightarrow real$
where *triangle-semiperimeter* $A\ B\ C =$
 $(dist\ B\ C + dist\ C\ A + dist\ A\ B) / 2$

definition *triangle-incenter* :: $real \wedge 2 \Rightarrow real \wedge 2 \Rightarrow real \wedge 2 \Rightarrow real \wedge 2$
where *triangle-incenter* $A\ B\ C =$
 $(let\ a = dist\ B\ C;\ b = dist\ C\ A;\ c = dist\ A\ B\ in$
 $inverse\ (a + b + c) *_{\mathbb{R}} (a *_{\mathbb{R}} A + b *_{\mathbb{R}} B + c *_{\mathbb{R}} C))$

definition *conway-point-past-first* :: $'a::real-normed-vector \Rightarrow 'a \Rightarrow real \Rightarrow 'a$
where *conway-point-past-first* $X\ Y\ d = X + (d / dist\ X\ Y) *_{\mathbb{R}} (X - Y)$

definition *conway-point-past-second* :: $'a::real-normed-vector \Rightarrow 'a \Rightarrow real \Rightarrow 'a$
where *conway-point-past-second* $X\ Y\ d = Y + (d / dist\ X\ Y) *_{\mathbb{R}} (Y - X)$

definition *side-contact-point* :: $'a::real-normed-vector \Rightarrow 'a \Rightarrow real \Rightarrow real \Rightarrow 'a$
where *side-contact-point* $X\ Y\ p\ q =$
 $(let\ L = dist\ X\ Y;\ s = (p + q + L) / 2\ in\ X + ((s - p) / L) *_{\mathbb{R}} (Y - X))$

definition *conway-A-c* :: $real \wedge 2 \Rightarrow real \wedge 2 \Rightarrow real \wedge 2 \Rightarrow real \wedge 2$
where *conway-A-c* $A\ B\ C = conway-point-past-first\ A\ B\ (dist\ B\ C)$

definition *conway-B-c* :: $\text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2$
where *conway-B-c* $A\ B\ C = \text{conway-point-past-second}\ A\ B\ (\text{dist}\ C\ A)$

definition *conway-A-b* :: $\text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2$
where *conway-A-b* $A\ B\ C = \text{conway-point-past-first}\ A\ C\ (\text{dist}\ B\ C)$

definition *conway-C-b* :: $\text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2$
where *conway-C-b* $A\ B\ C = \text{conway-point-past-second}\ A\ C\ (\text{dist}\ A\ B)$

definition *conway-B-a* :: $\text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2$
where *conway-B-a* $A\ B\ C = \text{conway-point-past-first}\ B\ C\ (\text{dist}\ C\ A)$

definition *conway-C-a* :: $\text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2$
where *conway-C-a* $A\ B\ C = \text{conway-point-past-second}\ B\ C\ (\text{dist}\ A\ B)$

definition *contact-AB* :: $\text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2$
where *contact-AB* $A\ B\ C = \text{side-contact-point}\ A\ B\ (\text{dist}\ B\ C)\ (\text{dist}\ C\ A)$

definition *contact-AC* :: $\text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2$
where *contact-AC* $A\ B\ C = \text{side-contact-point}\ A\ C\ (\text{dist}\ B\ C)\ (\text{dist}\ A\ B)$

definition *contact-BC* :: $\text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2$
where *contact-BC* $A\ B\ C = \text{side-contact-point}\ B\ C\ (\text{dist}\ C\ A)\ (\text{dist}\ A\ B)$

definition *triangle-inradius* :: $\text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real} \wedge 2 \Rightarrow \text{real}$
where *triangle-inradius* $A\ B\ C =$
 $\text{dist}\ (\text{triangle-incenter}\ A\ B\ C)\ (\text{contact-AB}\ A\ B\ C)$

lemma *noncollinear-triangle-distinct*:
assumes $\neg \text{collinear}\ \{A, B, C\}$
shows $A \neq B \wedge A \neq C \wedge B \neq C$
proof –
show $A \neq B$
proof
assume $A = B$
then have $\{A, B, C\} = \{A, C\}$
by *auto*
then have $\text{collinear}\ \{A, B, C\}$
by *simp*
with *assms* **show** *False*
by *contradiction*
qed
show $A \neq C$
proof
assume $A = C$
then have $\{A, B, C\} = \{B, C\}$
by *auto*
then have $\text{collinear}\ \{A, B, C\}$
by *simp*

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    with assms show False
      by contradiction
  qed
show  $B \neq C$ 
proof
  assume  $B = C$ 
  then have  $\{A, B, C\} = \{A, C\}$ 
    by auto
  then have collinear  $\{A, B, C\}$ 
    by simp
  with assms show False
    by contradiction
  qed
qed

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lemma inner-side-vectors:
  fixes  $X\ Y\ Z :: 'a::real-inner$ 
  shows  $(Z - X) \cdot (Y - X) =$ 
     $((dist\ Z\ X)^2 + (dist\ X\ Y)^2 - (dist\ Y\ Z)^2) / 2$ 
proof -
  have  $(Z - X) \cdot (Y - X) =$ 
     $((norm\ (Z - X))^2 + (norm\ (Y - X))^2 - (norm\ ((Z - X) - (Y - X)))^2)$ 
  / 2
    by (rule dot-norm-neg)
  also have  $\dots = ((dist\ Z\ X)^2 + (dist\ X\ Y)^2 - (dist\ Y\ Z)^2) / 2$ 
    by (simp add: dist-norm norm-minus-commute)
  finally show ?thesis .
qed

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lemma weighted-incenter-side-contact-orthogonal:
  fixes  $X\ Y\ Z :: 'a::real-inner$ 
  assumes  $X \neq Y$ 
  defines  $a \equiv dist\ Y\ Z$ 
  defines  $b \equiv dist\ Z\ X$ 
  defines  $c \equiv dist\ X\ Y$ 
  defines  $S \equiv a + b + c$ 
  defines  $I \equiv inverse\ S *_R (a *_R X + b *_R Y + c *_R Z)$ 
  defines  $T \equiv side-contact-point\ X\ Y\ a\ b$ 
  shows orthogonal  $(I - T) (Y - X)$ 
proof -
  have cpos:  $0 < c$ 
    using assms unfolding c-def by simp
  have a-nonneg:  $0 \leq a$ 
    unfolding a-def by simp
  have b-nonneg:  $0 \leq b$ 
    unfolding b-def by simp
  have Spos:  $0 < S$ 
    using a-nonneg b-nonneg cpos unfolding S-def by linarith
  define  $u$  where  $u = Y - X$ 

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define  $v$  where  $v = Z - X$ 
have incenter-translate:
   $I - X = \text{inverse } S *_R (b *_R u + c *_R v)$ 
proof -
  have numerator:
     $a *_R X + b *_R Y + c *_R Z =$ 
     $S *_R X + (b *_R (Y - X) + c *_R (Z - X))$ 
    unfolding S-def by (simp add: algebra-simps scaleR-diff-right)
  have  $I = X + \text{inverse } S *_R (b *_R (Y - X) + c *_R (Z - X))$ 
    unfolding I-def numerator
    using Spos by (simp add: scaleR-add-right)
  then show ?thesis
    unfolding u-def v-def by simp
qed
have contact-translate:
   $T - X = ((b + c - a) / (2 * c)) *_R u$ 
  unfolding T-def side-contact-point-def a-def b-def c-def Let-def
  by (simp add: field-simps u-def)
have side-square:  $u \cdot u = c^2$ 
proof -
  have  $(Y - X) \cdot (Y - X) = (\text{norm } (Y - X))^2$ 
    by (simp add: power2-norm-eq-inner)
  also have  $\dots = c^2$ 
    unfolding c-def by (simp add: dist-norm norm-minus-commute)
  finally show ?thesis
    unfolding u-def .
qed
have side-cross:  $v \cdot u = (b^2 + c^2 - a^2) / 2$ 
  unfolding u-def v-def a-def b-def c-def by (rule inner-side-vectors)
have  $(I - T) \cdot (Y - X) =$ 
   $((I - X) - (T - X)) \cdot u$ 
  unfolding u-def by simp
also have  $\dots =$ 
   $(\text{inverse } S *_R (b *_R u + c *_R v) -$ 
   $((b + c - a) / (2 * c)) *_R u) \cdot u$ 
  unfolding incenter-translate contact-translate by simp
also have  $\dots =$ 
   $\text{inverse } S * (b * c^2 + c * ((b^2 + c^2 - a^2) / 2)) -$ 
   $((b + c - a) / (2 * c)) * c^2$ 
  by (simp add: inner-simps side-square side-cross)
also have  $\dots = 0$ 
  using cpos Spots
  by (simp add: field-simps power2-eq-square S-def)
finally have  $(I - T) \cdot (Y - X) = 0$  .
then show ?thesis
  unfolding orthogonal-def by simp
qed

lemma incenter-contact-dist-sq-scalar:

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fixes  $a\ b\ c\ S :: \text{real}$ 
assumes  $c \neq 0$ 
and  $S \neq 0$ 
and  $S = a + b + c$ 
shows

$$\begin{aligned} & (\text{inverse } S * b - (b + c - a) / (2 * c))^2 * c^2 + \\ & (\text{inverse } S * c)^2 * b^2 + \\ & 2 * (\text{inverse } S * b - (b + c - a) / (2 * c)) * (\text{inverse } S * c) * \\ & ((b^2 + c^2 - a^2) / 2) = \\ & ((b + c - a) * (a + c - b) * (a + b - c)) / (4 * S) \end{aligned}$$

proof -
define  $D$  where  $D = a^2 - b^2 - c^2$ 
have alpha-eq:

$$\text{inverse } S * b - (b + c - a) / (2 * c) = D / (2 * c * S)$$

using assms unfolding  $D\text{-def}$ 
by (simp add: field-simps power2-eq-square)
have beta-eq:  $\text{inverse } S * c = c / S$ 
by (simp add: divide-inverse mult.commute)
have cross-eq:  $(b^2 + c^2 - a^2) / 2 = - D / 2$ 
unfolding  $D\text{-def}$  by simp
have

$$\begin{aligned} & (\text{inverse } S * b - (b + c - a) / (2 * c))^2 * c^2 + \\ & (\text{inverse } S * c)^2 * b^2 + \\ & 2 * (\text{inverse } S * b - (b + c - a) / (2 * c)) * (\text{inverse } S * c) * \\ & ((b^2 + c^2 - a^2) / 2) = \\ & (D / (2 * c * S))^2 * c^2 + (c / S)^2 * b^2 + \\ & 2 * (D / (2 * c * S)) * (c / S) * (- D / 2) \end{aligned}$$

by (simp add: alpha-eq beta-eq cross-eq)
also have  $\dots = (4 * b^2 * c^2 - D^2) / (4 * S^2)$ 
using assms(1,2) by (simp add: field-simps power2-eq-square)
also have  $\dots = ((2 * b * c - D) * (2 * b * c + D)) / (4 * S^2)$ 
by (simp add: power2-eq-square algebra-simps)
also have  $\dots =$ 

$$(((b + c)^2 - a^2) * (a^2 - (b - c)^2)) / (4 * S^2)$$

unfolding  $D\text{-def}$  by (simp add: power2-eq-square algebra-simps)
also have  $\dots =$ 

$$((S * (b + c - a)) * ((a + c - b) * (a + b - c))) / (4 * S^2)$$

using assms by (simp add: power2-eq-square algebra-simps)
also have  $\dots =$ 

$$((b + c - a) * (a + c - b) * (a + b - c)) / (4 * S)$$

using assms(2) by (simp add: field-simps power2-eq-square)
finally show ?thesis .
qed

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lemma weighted-incenter-side-contact-dist-sq:
fixes  $X\ Y\ Z :: 'a::\text{real-inner}$ 
assumes  $X \neq Y$ 
defines  $a \equiv \text{dist } Y\ Z$ 
defines  $b \equiv \text{dist } Z\ X$ 

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defines  $c \equiv \text{dist } X \ Y$ 
defines  $S \equiv a + b + c$ 
defines  $I \equiv \text{inverse } S *_R (a *_R X + b *_R Y + c *_R Z)$ 
defines  $T \equiv \text{side-contact-point } X \ Y \ a \ b$ 
shows  $(\text{dist } I \ T)^2 =$ 
   $((b + c - a) * (a + c - b) * (a + b - c)) / (4 * S)$ 
proof -
  have  $cpos: 0 < c$ 
    using assms unfolding c-def by simp
  have  $a\text{-nonneg}: 0 \leq a$ 
    unfolding a-def by simp
  have  $b\text{-nonneg}: 0 \leq b$ 
    unfolding b-def by simp
  have  $Spos: 0 < S$ 
    using a-nonneg b-nonneg cpos unfolding S-def by linarith
  define  $u$  where  $u = Y - X$ 
  define  $v$  where  $v = Z - X$ 
  have incenter-translate:
     $I - X = \text{inverse } S *_R (b *_R u + c *_R v)$ 
  proof -
    have numerator:
       $a *_R X + b *_R Y + c *_R Z =$ 
       $S *_R X + (b *_R (Y - X) + c *_R (Z - X))$ 
      unfolding S-def by (simp add: algebra-simps scaleR-diff-right)
    have  $I = X + \text{inverse } S *_R (b *_R (Y - X) + c *_R (Z - X))$ 
      unfolding I-def numerator
      using Spos by (simp add: scaleR-add-right)
    then show ?thesis
      unfolding u-def v-def by simp
  qed
  have contact-translate:
     $T - X = ((b + c - a) / (2 * c)) *_R u$ 
    unfolding T-def side-contact-point-def a-def b-def c-def Let-def
    by (simp add: field-simps u-def)
  have diff-translate:
     $I - T =$ 
     $\text{inverse } S *_R (b *_R u + c *_R v) -$ 
     $((b + c - a) / (2 * c)) *_R u$ 
  proof -
    have  $I - T = (I - X) - (T - X)$ 
      by simp
    also have  $\dots =$ 
     $\text{inverse } S *_R (b *_R u + c *_R v) -$ 
     $((b + c - a) / (2 * c)) *_R u$ 
    unfolding incenter-translate contact-translate ..
    finally show ?thesis .
  qed
  have side-square:  $u \cdot u = c^2$ 
  proof -

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have (Y - X) · (Y - X) = (norm (Y - X))2
  by (simp add: power2-norm-eq-inner)
also have ... = c2
  unfolding c-def by (simp add: dist-norm norm-minus-commute)
finally show ?thesis
  unfolding u-def .
qed
have vertex-square: v · v = b2
proof -
  have (Z - X) · (Z - X) = (norm (Z - X))2
    by (simp add: power2-norm-eq-inner)
  also have ... = b2
    unfolding b-def by (simp add: dist-norm norm-minus-commute)
  finally show ?thesis
    unfolding v-def .
qed
have side-cross: v · u = (b2 + c2 - a2) / 2
  unfolding u-def v-def a-def b-def c-def by (rule inner-side-vectors)
have side-cross': u · v = (b2 + c2 - a2) / 2
  using side-cross by (simp add: inner-commute)
define alpha where alpha = inverse S * b - ((b + c - a) / (2 * c))
define beta where beta = inverse S * c
have linear-diff:
  inverse S *R (b *R u + c *R v) -
  ((b + c - a) / (2 * c)) *R u =
  alpha *R u + beta *R v
  unfolding alpha-def beta-def
  by (simp add: algebra-simps scaleR-add-right scaleR-diff-left)
have (dist I T)2 = (I - T) · (I - T)
  by (simp add: dist-norm power2-norm-eq-inner)
also have ... =
  (inverse S *R (b *R u + c *R v) -
  ((b + c - a) / (2 * c)) *R u) ·
  (inverse S *R (b *R u + c *R v) -
  ((b + c - a) / (2 * c)) *R u)
  unfolding diff-translate ..
also have ... =
  (alpha *R u + beta *R v) · (alpha *R u + beta *R v)
  unfolding linear-diff ..
also have ... =
  alpha2 * c2 + beta2 * b2 +
  2 * alpha * beta * ((b2 + c2 - a2) / 2)
  by (simp add: inner-simps side-square vertex-square side-cross side-cross'
  power2-eq-square algebra-simps) algebra
also have ... =
  (inverse S * b - ((b + c - a) / (2 * c)))2 * c2 +
  (inverse S * c)2 * b2 +
  2 * (inverse S * b - ((b + c - a) / (2 * c))) *
  (inverse S * c) * ((b2 + c2 - a2) / 2)

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    unfolding alpha-def beta-def ..
  also have ... =
    ((b + c - a) * (a + c - b) * (a + b - c)) / (4 * S)
  proof (rule incenter-contact-dist-sq-scalar)
    show c ≠ 0
      using cpos by simp
    show S ≠ 0
      using Spos by simp
    show S = a + b + c
      unfolding S-def ..
  qed
  finally show ?thesis .
qed

```

```

lemma triangle-incenter-contact-AB-orthogonal:
  fixes A B C :: real ^ 2
  assumes A ≠ B
  shows orthogonal (triangle-incenter A B C - contact-AB A B C) (B - A)
  using weighted-incenter-side-contact-orthogonal[OF assms, of C]
  unfolding triangle-incenter-def contact-AB-def Let-def
  by simp

```

```

lemma triangle-incenter-contact-AC-orthogonal:
  fixes A B C :: real ^ 2
  assumes A ≠ C
  shows orthogonal (triangle-incenter A B C - contact-AC A B C) (C - A)
  using weighted-incenter-side-contact-orthogonal[OF assms, of B]
  unfolding triangle-incenter-def contact-AC-def Let-def
  by (simp add: algebra-simps dist-commute)

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```

lemma triangle-incenter-contact-BC-orthogonal:
  fixes A B C :: real ^ 2
  assumes B ≠ C
  shows orthogonal (triangle-incenter A B C - contact-BC A B C) (C - B)
  using weighted-incenter-side-contact-orthogonal[OF assms, of A]
  unfolding triangle-incenter-def contact-BC-def Let-def
  by (simp add: algebra-simps dist-commute)

```

```

lemma triangle-inradius-nonneg [simp]: 0 ≤ triangle-inradius A B C
  unfolding triangle-inradius-def by simp

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lemma triangle-incenter-contact-AB-dist:
  dist (triangle-incenter A B C) (contact-AB A B C) = triangle-inradius A B C
  unfolding triangle-inradius-def ..

```

```

lemma triangle-incenter-contact-AC-dist:
  fixes A B C :: real ^ 2
  assumes triangle: ¬ collinear {A, B, C}
  shows dist (triangle-incenter A B C) (contact-AC A B C) = triangle-inradius A

```

```

B C
proof -
  let ?I = triangle-incenter A B C
  let ?a = dist B C
  let ?b = dist C A
  let ?c = dist A B
  let ?H = ((?b + ?c - ?a) * (?a + ?c - ?b) * (?a + ?b - ?c)) /
    (4 * (?a + ?b + ?c))
  have sides: A ≠ B A ≠ C
    using noncollinear-triangle-distinct[OF triangle] by blast+
  have AB-sq: (dist ?I (contact-AB A B C))2 = ?H
    using weighted-incenter-side-contact-dist-sq[OF sides(1), of C]
    unfolding triangle-incenter-def contact-AB-def Let-def
    by (simp add: dist-commute algebra-simps)
  have AC-sq: (dist ?I (contact-AC A B C))2 = ?H
    using weighted-incenter-side-contact-dist-sq[OF sides(2), of B]
    unfolding triangle-incenter-def contact-AC-def Let-def
    by (simp add: dist-commute algebra-simps)
  show ?thesis
    unfolding triangle-inradius-def
  proof (rule power2-eq-imp-eq)
    show (dist ?I (contact-AC A B C))2 =
      (dist ?I (contact-AB A B C))2
      using AC-sq AB-sq by simp
    show 0 ≤ dist ?I (contact-AC A B C)
      by simp
    show 0 ≤ dist ?I (contact-AB A B C)
      by simp
  qed
qed

lemma triangle-incenter-contact-BC-dist:
  fixes A B C :: real ^ 2
  assumes triangle: ¬ collinear {A, B, C}
  shows dist (triangle-incenter A B C) (contact-BC A B C) = triangle-inradius A
B C
proof -
  let ?I = triangle-incenter A B C
  let ?a = dist B C
  let ?b = dist C A
  let ?c = dist A B
  let ?H = ((?b + ?c - ?a) * (?a + ?c - ?b) * (?a + ?b - ?c)) /
    (4 * (?a + ?b + ?c))
  have sides: A ≠ B B ≠ C
    using noncollinear-triangle-distinct[OF triangle] by blast+
  have AB-sq: (dist ?I (contact-AB A B C))2 = ?H
    using weighted-incenter-side-contact-dist-sq[OF sides(1), of C]
    unfolding triangle-incenter-def contact-AB-def Let-def
    by (simp add: dist-commute algebra-simps)

```

```

have BC-sq: (dist ?I (contact-BC A B C))2 = ?H
  using weighted-incenter-side-contact-dist-sq[OF sides(2), of A]
  unfolding triangle-incenter-def contact-BC-def Let-def
  by (simp add: dist-commute algebra-simps)
show ?thesis
  unfolding triangle-inradius-def
proof (rule power2-eq-imp-eq)
  show (dist ?I (contact-BC A B C))2 =
    (dist ?I (contact-AB A B C))2
    using BC-sq AB-sq by simp
  show 0 ≤ dist ?I (contact-BC A B C)
    by simp
  show 0 ≤ dist ?I (contact-AB A B C)
    by simp
qed
qed

lemma triangle-incenter-contacts-dist:
  fixes A B C :: real ^ 2
  assumes ¬ collinear {A, B, C}
  shows
    dist (triangle-incenter A B C) (contact-AB A B C) = triangle-inradius A B C
    dist (triangle-incenter A B C) (contact-AC A B C) = triangle-inradius A B C
    dist (triangle-incenter A B C) (contact-BC A B C) = triangle-inradius A B C
proof -
  show dist (triangle-incenter A B C) (contact-AB A B C) = triangle-inradius A
    B C
    by (rule triangle-incenter-contact-AB-dist)
  show dist (triangle-incenter A B C) (contact-AC A B C) = triangle-inradius A
    B C
    by (rule triangle-incenter-contact-AC-dist[OF assms])
  show dist (triangle-incenter A B C) (contact-BC A B C) = triangle-inradius A
    B C
    by (rule triangle-incenter-contact-BC-dist[OF assms])
qed

lemma triangle-incenter-contacts-orthogonal:
  fixes A B C :: real ^ 2
  assumes ¬ collinear {A, B, C}
  shows
    orthogonal (triangle-incenter A B C - contact-AB A B C) (B - A)
    orthogonal (triangle-incenter A B C - contact-AC A B C) (C - A)
    orthogonal (triangle-incenter A B C - contact-BC A B C) (C - B)
proof -
  have sides: A ≠ B A ≠ C B ≠ C
    using noncollinear-triangle-distinct[OF assms] by blast+
  show orthogonal (triangle-incenter A B C - contact-AB A B C) (B - A)
    by (rule triangle-incenter-contact-AB-orthogonal[OF sides(1)])
  show orthogonal (triangle-incenter A B C - contact-AC A B C) (C - A)

```

by (rule triangle-incenter-contact-AC-orthogonal[OF sides(2)])
 show orthogonal (triangle-incenter A B C - contact-BC A B C) (C - B)
 by (rule triangle-incenter-contact-BC-orthogonal[OF sides(3)])
 qed

lemma *dist-conway-point-past-first*:
 fixes X Y :: 'a::real-normed-vector
 assumes X ≠ Y 0 ≤ d
 shows dist X (conway-point-past-first X Y d) = d
proof -
 have 0 < dist X Y
 using assms by simp
 then show ?thesis
 using assms
 unfolding conway-point-past-first-def
 by (simp add: dist-norm norm-minus-commute)
 qed

lemma *dist-conway-point-past-second*:
 fixes X Y :: 'a::real-normed-vector
 assumes X ≠ Y 0 ≤ d
 shows dist Y (conway-point-past-second X Y d) = d
proof -
 have 0 < dist X Y
 using assms by simp
 then show ?thesis
 using assms
 unfolding conway-point-past-second-def
 by (simp add: dist-norm norm-minus-commute)
 qed

lemma *conway-point-past-first-contact*:
 fixes X Y :: 'a::real-normed-vector
 assumes X ≠ Y
 shows conway-point-past-first X Y p =
 side-contact-point X Y p q - (((p + q + dist X Y) / 2) / dist X Y) *_R (Y - X)
proof -
 let ?L = dist X Y
 let ?s = (p + q + ?L) / 2
 have Lpos: 0 < ?L
 using assms by simp
 have scalar: ((?s - p) / ?L - ?s / ?L) = - (p / ?L)
 using Lpos by (simp add: field-simps)
 have side-contact-point X Y p q - (?s / ?L) *_R (Y - X) =
 X + ((?s - p) / ?L) *_R (Y - X) - (?s / ?L) *_R (Y - X)
 unfolding side-contact-point-def Let-def by simp
 also have ... = X + (((?s - p) / ?L - ?s / ?L) *_R (Y - X))
 by (simp add: scaleR-diff-left [symmetric])

```

also have ... = X + (-(p / ?L) *R (Y - X))
  using scalar by simp
also have ... = X + (p / ?L) *R (X - Y)
  by (simp add: scaleR-right-diff-distrib)
also have ... = conway-point-past-first X Y p
  unfolding conway-point-past-first-def ..
finally show ?thesis
  by (rule sym)
qed

```

lemma *conway-point-past-second-contact*:

```

fixes X Y :: 'a::real-normed-vector
assumes X ≠ Y
shows conway-point-past-second X Y q =
  side-contact-point X Y p q + (((p + q + dist X Y) / 2) / dist X Y) *R (Y -
X)
proof -
  let ?L = dist X Y
  let ?s = (p + q + ?L) / 2
  have Lpos: 0 < ?L
    using assms by simp
  have scalar: ((?s - p) / ?L + ?s / ?L) = 1 + q / ?L
    using Lpos by (simp add: field-simps)
  have side-contact-point X Y p q + (?s / ?L) *R (Y - X) =
    X + ((?s - p) / ?L) *R (Y - X) + (?s / ?L) *R (Y - X)
    unfolding side-contact-point-def Let-def by simp
  also have ... = X + (((?s - p) / ?L + ?s / ?L) *R (Y - X))
    by (simp add: scaleR-left-distrib [symmetric])
  also have ... = X + (1 + q / ?L) *R (Y - X)
    using scalar by simp
  also have ... = Y + (q / ?L) *R (Y - X)
    by (simp add: scaleR-left-distrib)
  also have ... = conway-point-past-second X Y q
    unfolding conway-point-past-second-def ..
  finally show ?thesis
    by (rule sym)
qed

```

lemma *dist-side-contact-first*:

```

fixes X Y :: 'a::real-normed-vector
assumes X ≠ Y 0 ≤ p 0 ≤ q
shows dist (side-contact-point X Y p q) (conway-point-past-first X Y p) =
  (p + q + dist X Y) / 2
proof -
  let ?s = (p + q + dist X Y) / 2
  have Lpos: 0 < dist X Y
    using assms by simp
  have spos: 0 ≤ ?s
    using assms Lpos by simp

```

```

have point-eq:
  conway-point-past-first  $X\ Y\ p =$ 
    side-contact-point  $X\ Y\ p\ q - (?s / \text{dist } X\ Y) *_R (Y - X)$ 
  by (rule conway-point-past-first-contact[ $OF\ \text{assms}(1)$ ])
have diff-eq:
  side-contact-point  $X\ Y\ p\ q - \text{conway-point-past-first } X\ Y\ p =$ 
     $(?s / \text{dist } X\ Y) *_R (Y - X)$ 
  using point-eq by simp
have dist (side-contact-point  $X\ Y\ p\ q$ ) (conway-point-past-first  $X\ Y\ p$ ) =
  norm  $((?s / \text{dist } X\ Y) *_R (Y - X))$ 
  by (simp add: dist-norm diff-eq)
also have ... = abs  $(?s / \text{dist } X\ Y) * \text{norm } (Y - X)$ 
  by simp
also have ... =  $(?s / \text{dist } X\ Y) * \text{dist } X\ Y$ 
  using Lpos spos by (simp add: dist-norm norm-minus-commute)
also have ... = ?s
  using Lpos by simp
finally show ?thesis .
qed

lemma dist-side-contact-second:
  fixes  $X\ Y :: 'a::\text{real-normed-vector}$ 
  assumes  $X \neq Y\ 0 \leq p\ 0 \leq q$ 
  shows dist (side-contact-point  $X\ Y\ p\ q$ ) (conway-point-past-second  $X\ Y\ q$ ) =
     $(p + q + \text{dist } X\ Y) / 2$ 
proof -
  let ?s =  $(p + q + \text{dist } X\ Y) / 2$ 
  have Lpos:  $0 < \text{dist } X\ Y$ 
  using assms by simp
  have spos:  $0 \leq ?s$ 
  using assms Lpos by simp
  have point-eq:
    conway-point-past-second  $X\ Y\ q =$ 
      side-contact-point  $X\ Y\ p\ q + (?s / \text{dist } X\ Y) *_R (Y - X)$ 
  by (rule conway-point-past-second-contact[ $OF\ \text{assms}(1)$ ])
  have diff-eq:
    side-contact-point  $X\ Y\ p\ q - \text{conway-point-past-second } X\ Y\ q =$ 
       $-(?s / \text{dist } X\ Y) *_R (Y - X)$ 
  using point-eq by simp
  have dist (side-contact-point  $X\ Y\ p\ q$ ) (conway-point-past-second  $X\ Y\ q$ ) =
    norm  $((?s / \text{dist } X\ Y) *_R (Y - X))$ 
  by (simp add: dist-norm diff-eq)
  also have ... = abs  $(?s / \text{dist } X\ Y) * \text{norm } (Y - X)$ 
  by simp
  also have ... =  $(?s / \text{dist } X\ Y) * \text{dist } X\ Y$ 
  using Lpos spos by (simp add: dist-norm norm-minus-commute)
  also have ... = ?s
  using Lpos by simp
  finally show ?thesis .

```

qed

lemma *dist-side-contact-first'*:

fixes $X\ Y :: 'a::\text{real-normed-vector}$

assumes $X \neq Y\ 0 \leq p\ 0 \leq q$

shows $\text{dist} (\text{conway-point-past-first } X\ Y\ p) (\text{side-contact-point } X\ Y\ p\ q) =$
 $(p + q + \text{dist } X\ Y) / 2$

using *dist-side-contact-first*[*OF assms*] **by** (*simp add: dist-commute*)

lemma *dist-side-contact-second'*:

fixes $X\ Y :: 'a::\text{real-normed-vector}$

assumes $X \neq Y\ 0 \leq p\ 0 \leq q$

shows $\text{dist} (\text{conway-point-past-second } X\ Y\ q) (\text{side-contact-point } X\ Y\ p\ q) =$
 $(p + q + \text{dist } X\ Y) / 2$

using *dist-side-contact-second*[*OF assms*] **by** (*simp add: dist-commute*)

lemma *orthogonal-conway-point-past-first-contact*:

fixes $X\ Y\ v :: 'a::\text{real-inner}$

assumes $X \neq Y\ \text{orthogonal } v\ (Y - X)$

shows $\text{orthogonal } v\ (\text{conway-point-past-first } X\ Y\ p - \text{side-contact-point } X\ Y\ p\ q)$

proof –

let $?s = (p + q + \text{dist } X\ Y) / 2$

have *point-eq*:

$\text{conway-point-past-first } X\ Y\ p - \text{side-contact-point } X\ Y\ p\ q =$
 $-(?s / \text{dist } X\ Y) *_R (Y - X)$

using *conway-point-past-first-contact*[*OF assms(1)*, **where** $p=p$ **and** $q=q$] **by**

simp

show *?thesis*

using *assms(2)* **unfolding** *point-eq* *orthogonal-def* **by** *simp*

qed

lemma *orthogonal-conway-point-past-second-contact*:

fixes $X\ Y\ v :: 'a::\text{real-inner}$

assumes $X \neq Y\ \text{orthogonal } v\ (Y - X)$

shows $\text{orthogonal } v\ (\text{conway-point-past-second } X\ Y\ q - \text{side-contact-point } X\ Y\ p\ q)$

proof –

let $?s = (p + q + \text{dist } X\ Y) / 2$

have *point-eq*:

$\text{conway-point-past-second } X\ Y\ q - \text{side-contact-point } X\ Y\ p\ q =$
 $(?s / \text{dist } X\ Y) *_R (Y - X)$

using *conway-point-past-second-contact*[*OF assms(1)*, **where** $p=p$ **and** $q=q$]

by *simp*

show *?thesis*

using *assms(2)* **unfolding** *point-eq* *orthogonal-def* **by** *simp*

qed

lemma *dist-from-orthogonal-foot*:

fixes $I\ T\ P :: 'a::\text{real-inner}$


```

assumes orth: orthogonal ( $I - T$ ) ( $P - T$ )
  and IT: dist  $I\ T = r$ 
  and TP: dist  $T\ P = s$ 
  and r-nonneg:  $0 \leq r$ 
  and s-nonneg:  $0 \leq s$ 
shows dist  $I\ P = \text{sqrt } (r^2 + s^2)$ 
proof -
  have inner-flip:  $(I - T) \cdot (T - P) = - ((I - T) \cdot (P - T))$ 
    by (simp add: inner-diff-right)
  have orth': orthogonal ( $I - T$ ) ( $T - P$ )
    using orth inner-flip by (simp add: orthogonal-def)
  have Pythagoras:
     $(\text{norm } ((I - T) + (T - P)))^2 =$ 
     $(\text{norm } (I - T))^2 + (\text{norm } (T - P))^2$ 
    by (rule norm-add-Pythagorean[OF orth'])
  have square-dist:
     $(\text{dist } I\ P)^2 = r^2 + s^2$ 
    using Pythagoras IT TP
    by (simp add: dist-norm norm-minus-commute)
  have dist I P = sqrt ((dist I P)^2)
    by simp
  also have  $\dots = \text{sqrt } (r^2 + s^2)$ 
    using square-dist by simp
  finally show ?thesis .
qed

lemma conway-points-have-declared-lengths:
assumes  $A \neq B\ A \neq C\ B \neq C$ 
shows
  dist  $A\ (\text{conway-A-c } A\ B\ C) = \text{dist } B\ C$ 
  dist  $B\ (\text{conway-B-c } A\ B\ C) = \text{dist } C\ A$ 
  dist  $A\ (\text{conway-A-b } A\ B\ C) = \text{dist } B\ C$ 
  dist  $C\ (\text{conway-C-b } A\ B\ C) = \text{dist } A\ B$ 
  dist  $B\ (\text{conway-B-a } A\ B\ C) = \text{dist } C\ A$ 
  dist  $C\ (\text{conway-C-a } A\ B\ C) = \text{dist } A\ B$ 
using assms
unfolding conway-A-c-def conway-B-c-def conway-A-b-def conway-C-b-def
  conway-B-a-def conway-C-a-def
by (simp-all add: dist-conway-point-past-first dist-conway-point-past-second)

lemma conway-points-about-contacts:
assumes  $A \neq B\ A \neq C\ B \neq C$ 
defines  $s \equiv \text{triangle-semiperimeter } A\ B\ C$ 
shows
  dist (contact-AB  $A\ B\ C$ ) (conway-A-c  $A\ B\ C$ ) =  $s$ 
  dist (contact-AB  $A\ B\ C$ ) (conway-B-c  $A\ B\ C$ ) =  $s$ 
  dist (contact-AC  $A\ B\ C$ ) (conway-A-b  $A\ B\ C$ ) =  $s$ 
  dist (contact-AC  $A\ B\ C$ ) (conway-C-b  $A\ B\ C$ ) =  $s$ 
  dist (contact-BC  $A\ B\ C$ ) (conway-B-a  $A\ B\ C$ ) =  $s$ 

```

$\text{dist } (\text{contact-BC } A \ B \ C) \ (\text{conway-C-a } A \ B \ C) = s$
using *assms*
unfolding *s-def triangle-semiperimeter-def*
 $\text{contact-AB-def contact-AC-def contact-BC-def}$
 $\text{conway-A-c-def conway-B-c-def conway-A-b-def conway-C-b-def}$
 $\text{conway-B-a-def conway-C-a-def}$
by (*simp-all add: dist-side-contact-first dist-side-contact-second*
 $\text{dist-side-contact-first' dist-side-contact-second' dist-commute}$)

lemma *conway-points-on-circle-from-contact-feet:*
fixes $A \ B \ C \ I :: \text{real} \wedge 2$
assumes *sides:* $A \neq B \ A \neq C \ B \neq C$
and *r-nonneg:* $0 \leq r$
and *foot-AB:* $\text{dist } I \ (\text{contact-AB } A \ B \ C) = r$
 $\text{orthogonal } (I - \text{contact-AB } A \ B \ C) \ (B - A)$
and *foot-AC:* $\text{dist } I \ (\text{contact-AC } A \ B \ C) = r$
 $\text{orthogonal } (I - \text{contact-AC } A \ B \ C) \ (C - A)$
and *foot-BC:* $\text{dist } I \ (\text{contact-BC } A \ B \ C) = r$
 $\text{orthogonal } (I - \text{contact-BC } A \ B \ C) \ (C - B)$
defines $s \equiv \text{triangle-semiperimeter } A \ B \ C$
shows
 $\text{dist } I \ (\text{conway-A-c } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I \ (\text{conway-B-c } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I \ (\text{conway-A-b } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I \ (\text{conway-C-b } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I \ (\text{conway-B-a } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I \ (\text{conway-C-a } A \ B \ C) = \text{sqrt } (r^2 + s^2)$

proof –
have *s-nonneg:* $0 \leq s$
unfolding *s-def triangle-semiperimeter-def* **by** *simp*
have *contact-dists:*
 $\text{dist } (\text{contact-AB } A \ B \ C) \ (\text{conway-A-c } A \ B \ C) = s$
 $\text{dist } (\text{contact-AB } A \ B \ C) \ (\text{conway-B-c } A \ B \ C) = s$
 $\text{dist } (\text{contact-AC } A \ B \ C) \ (\text{conway-A-b } A \ B \ C) = s$
 $\text{dist } (\text{contact-AC } A \ B \ C) \ (\text{conway-C-b } A \ B \ C) = s$
 $\text{dist } (\text{contact-BC } A \ B \ C) \ (\text{conway-B-a } A \ B \ C) = s$
 $\text{dist } (\text{contact-BC } A \ B \ C) \ (\text{conway-C-a } A \ B \ C) = s$
using *conway-points-about-contacts[OF sides, folded s-def]* .
have *orth-Ac:*
 $\text{orthogonal } (I - \text{contact-AB } A \ B \ C) \ (\text{conway-A-c } A \ B \ C - \text{contact-AB } A \ B \ C)$
using *foot-AB(2)*
unfolding *conway-A-c-def contact-AB-def*
by (*rule orthogonal-conway-point-past-first-contact[OF sides(1)]*)
have *orth-Bc:*
 $\text{orthogonal } (I - \text{contact-AB } A \ B \ C) \ (\text{conway-B-c } A \ B \ C - \text{contact-AB } A \ B \ C)$
using *foot-AB(2)*
unfolding *conway-B-c-def contact-AB-def*

by (rule *orthogonal-conway-point-past-second-contact*[*OF sides*(1)])
have *orth-Ab*:
orthogonal (*I* – *contact-AC A B C*) (*conway-A-b A B C* – *contact-AC A B C*)
using *foot-AC*(2)
unfolding *conway-A-b-def contact-AC-def*
by (rule *orthogonal-conway-point-past-first-contact*[*OF sides*(2)])
have *orth-Cb*:
orthogonal (*I* – *contact-AC A B C*) (*conway-C-b A B C* – *contact-AC A B C*)
using *foot-AC*(2)
unfolding *conway-C-b-def contact-AC-def*
by (rule *orthogonal-conway-point-past-second-contact*[*OF sides*(2)])
have *orth-Ba*:
orthogonal (*I* – *contact-BC A B C*) (*conway-B-a A B C* – *contact-BC A B*
C)
using *foot-BC*(2)
unfolding *conway-B-a-def contact-BC-def*
by (rule *orthogonal-conway-point-past-first-contact*[*OF sides*(3)])
have *orth-Ca*:
orthogonal (*I* – *contact-BC A B C*) (*conway-C-a A B C* – *contact-BC A B*
C)
using *foot-BC*(2)
unfolding *conway-C-a-def contact-BC-def*
by (rule *orthogonal-conway-point-past-second-contact*[*OF sides*(3)])
show *dist I (conway-A-c A B C) = sqrt (r² + s²)*
by (rule *dist-from-orthogonal-foot*[*OF orth-Ac foot-AB*(1) *contact-dists*(1) *r-nonneg*
s-nonneg])
show *dist I (conway-B-c A B C) = sqrt (r² + s²)*
by (rule *dist-from-orthogonal-foot*[*OF orth-Bc foot-AB*(1) *contact-dists*(2) *r-nonneg*
s-nonneg])
show *dist I (conway-A-b A B C) = sqrt (r² + s²)*
by (rule *dist-from-orthogonal-foot*[*OF orth-Ab foot-AC*(1) *contact-dists*(3) *r-nonneg*
s-nonneg])
show *dist I (conway-C-b A B C) = sqrt (r² + s²)*
by (rule *dist-from-orthogonal-foot*[*OF orth-Cb foot-AC*(1) *contact-dists*(4) *r-nonneg*
s-nonneg])
show *dist I (conway-B-a A B C) = sqrt (r² + s²)*
by (rule *dist-from-orthogonal-foot*[*OF orth-Ba foot-BC*(1) *contact-dists*(5) *r-nonneg*
s-nonneg])
show *dist I (conway-C-a A B C) = sqrt (r² + s²)*
by (rule *dist-from-orthogonal-foot*[*OF orth-Ca foot-BC*(1) *contact-dists*(6) *r-nonneg*
s-nonneg])
qed

lemma *conway-points-on-circle-from-noncollinear-contact-feet*:

fixes *A B C I* :: *real* ^ 2
assumes *triangle*: $\neg \text{collinear } \{A, B, C\}$
and *r-nonneg*: $0 \leq r$
and *foot-AB*: *dist I (contact-AB A B C) = r*
orthogonal (I – contact-AB A B C) (B – A)

and *foot-AC*: $\text{dist } I (\text{contact-AC } A \ B \ C) = r$
 $\text{orthogonal } (I - \text{contact-AC } A \ B \ C) (C - A)$
and *foot-BC*: $\text{dist } I (\text{contact-BC } A \ B \ C) = r$
 $\text{orthogonal } (I - \text{contact-BC } A \ B \ C) (C - B)$
defines $s \equiv \text{triangle-semiperimeter } A \ B \ C$
shows
 $\text{dist } I (\text{conway-A-c } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I (\text{conway-B-c } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I (\text{conway-A-b } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I (\text{conway-C-b } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I (\text{conway-B-a } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I (\text{conway-C-a } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
proof –
have *sides*: $A \neq B \ A \neq C \ B \neq C$
using *noncollinear-triangle-distinct*[*OF triangle*] **by** *blast+*
show $\text{dist } I (\text{conway-A-c } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
and $\text{dist } I (\text{conway-B-c } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
and $\text{dist } I (\text{conway-A-b } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
and $\text{dist } I (\text{conway-C-b } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
and $\text{dist } I (\text{conway-B-a } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
and $\text{dist } I (\text{conway-C-a } A \ B \ C) = \text{sqrt } (r^2 + s^2)$
using *conway-points-on-circle-from-contact-feet*
 $[OF \text{ sides } r\text{-nonneg } \text{foot-AB } \text{foot-AC } \text{foot-BC}, \text{ folded } s\text{-def}]$.
qed

theorem *conway-circle-theorem*:
fixes $A \ B \ C :: \text{real}^{\wedge} 2$
assumes *triangle*: $\neg \text{collinear } \{A, B, C\}$
defines $I \equiv \text{triangle-incenter } A \ B \ C$
defines $r \equiv \text{triangle-inradius } A \ B \ C$
defines $s \equiv \text{triangle-semiperimeter } A \ B \ C$
defines $A\text{-c} \equiv \text{conway-A-c } A \ B \ C$
defines $B\text{-c} \equiv \text{conway-B-c } A \ B \ C$
defines $A\text{-b} \equiv \text{conway-A-b } A \ B \ C$
defines $C\text{-b} \equiv \text{conway-C-b } A \ B \ C$
defines $B\text{-a} \equiv \text{conway-B-a } A \ B \ C$
defines $C\text{-a} \equiv \text{conway-C-a } A \ B \ C$
shows
 $\text{dist } I \ A\text{-c} = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I \ B\text{-c} = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I \ A\text{-b} = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I \ C\text{-b} = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I \ B\text{-a} = \text{sqrt } (r^2 + s^2)$
 $\text{dist } I \ C\text{-a} = \text{sqrt } (r^2 + s^2)$
proof –
have *foot-dists*:
 $\text{dist } I (\text{contact-AB } A \ B \ C) = r$
 $\text{dist } I (\text{contact-AC } A \ B \ C) = r$
 $\text{dist } I (\text{contact-BC } A \ B \ C) = r$

```

using triangle-incenter-contacts-dist[OF triangle]
unfolding I-def r-def by simp-all
have foot-orthogonal:
  orthogonal (I - contact-AB A B C) (B - A)
  orthogonal (I - contact-AC A B C) (C - A)
  orthogonal (I - contact-BC A B C) (C - B)
using triangle-incenter-contacts-orthogonal[OF triangle]
unfolding I-def by simp-all
have r-nonneg:  $0 \leq r$ 
unfolding r-def by simp
have circle:
  dist I (conway-A-c A B C) = sqrt ( $r^2 + s^2$ )
  dist I (conway-B-c A B C) = sqrt ( $r^2 + s^2$ )
  dist I (conway-A-b A B C) = sqrt ( $r^2 + s^2$ )
  dist I (conway-C-b A B C) = sqrt ( $r^2 + s^2$ )
  dist I (conway-B-a A B C) = sqrt ( $r^2 + s^2$ )
  dist I (conway-C-a A B C) = sqrt ( $r^2 + s^2$ )
using conway-points-on-circle-from-noncollinear-contact-feet
  [OF triangle r-nonneg foot-dists(1) foot-orthogonal(1)
   foot-dists(2) foot-orthogonal(2) foot-dists(3) foot-orthogonal(3),
   folded s-def] .
show dist I A-c = sqrt ( $r^2 + s^2$ )
unfolding A-c-def by (rule circle(1))
show dist I B-c = sqrt ( $r^2 + s^2$ )
unfolding B-c-def by (rule circle(2))
show dist I A-b = sqrt ( $r^2 + s^2$ )
unfolding A-b-def by (rule circle(3))
show dist I C-b = sqrt ( $r^2 + s^2$ )
unfolding C-b-def by (rule circle(4))
show dist I B-a = sqrt ( $r^2 + s^2$ )
unfolding B-a-def by (rule circle(5))
show dist I C-a = sqrt ( $r^2 + s^2$ )
unfolding C-a-def by (rule circle(6))
qed

```

4.1 Wrapper formulations

The headline theorem above proves six separate distance equations. For downstream use we package the result in two equivalent forms: as a subset of an explicit *sphere* centred at the incenter, and as a single bounded-universal statement over the explicit six-point set. We also expose the construction correctness as a separate theorem: each Conway point lies on the intended side line and is at the declared distance from the adjacent vertex.

```

lemma conway-point-past-first-on-line:
  fixes X Y :: 'a::real-normed-vector
  shows conway-point-past-first X Y d = X + (- (d / dist X Y)) *R (Y - X)
  unfolding conway-point-past-first-def
  by (simp add: scaleR-minus-right algebra-simps)

```

lemma *conway-point-past-second-on-line*:
fixes $X\ Y :: 'a::\text{real-normed-vector}$
shows *conway-point-past-second* $X\ Y\ d$
 $= X + (1 + d / \text{dist}\ X\ Y) *_{\mathbb{R}} (Y - X)$
unfolding *conway-point-past-second-def*
by (*simp add: scaleR-left-distrib*)

theorem *conway-construction-correct*:
fixes $A\ B\ C :: \text{real}^2$
assumes *triangle*: $\neg \text{collinear}\ \{A, B, C\}$
shows
— The six Conway points lie on the three side lines of the triangle.
 $\exists t::\text{real. } \text{conway-A-c}\ A\ B\ C = A + t *_{\mathbb{R}} (B - A)$
 $\exists t::\text{real. } \text{conway-B-c}\ A\ B\ C = A + t *_{\mathbb{R}} (B - A)$
 $\exists t::\text{real. } \text{conway-A-b}\ A\ B\ C = A + t *_{\mathbb{R}} (C - A)$
 $\exists t::\text{real. } \text{conway-C-b}\ A\ B\ C = A + t *_{\mathbb{R}} (C - A)$
 $\exists t::\text{real. } \text{conway-B-a}\ A\ B\ C = B + t *_{\mathbb{R}} (C - B)$
 $\exists t::\text{real. } \text{conway-C-a}\ A\ B\ C = B + t *_{\mathbb{R}} (C - B)$
— Their distances from the adjacent vertex equal the opposite side length.
 $\text{dist}\ A\ (\text{conway-A-c}\ A\ B\ C) = \text{dist}\ B\ C$
 $\text{dist}\ B\ (\text{conway-B-c}\ A\ B\ C) = \text{dist}\ C\ A$
 $\text{dist}\ A\ (\text{conway-A-b}\ A\ B\ C) = \text{dist}\ B\ C$
 $\text{dist}\ C\ (\text{conway-C-b}\ A\ B\ C) = \text{dist}\ A\ B$
 $\text{dist}\ B\ (\text{conway-B-a}\ A\ B\ C) = \text{dist}\ C\ A$
 $\text{dist}\ C\ (\text{conway-C-a}\ A\ B\ C) = \text{dist}\ A\ B$

proof —
have *sides*: $A \neq B\ A \neq C\ B \neq C$
using *noncollinear-triangle-distinct[OF triangle]* **by** *blast+*
show
 $\exists t::\text{real. } \text{conway-A-c}\ A\ B\ C = A + t *_{\mathbb{R}} (B - A)$
unfolding *conway-A-c-def*
using *conway-point-past-first-on-line[of A B dist B C]* **by** *blast*
show
 $\exists t::\text{real. } \text{conway-B-c}\ A\ B\ C = A + t *_{\mathbb{R}} (B - A)$
unfolding *conway-B-c-def*
using *conway-point-past-second-on-line[of A B dist C A]* **by** *blast*
show
 $\exists t::\text{real. } \text{conway-A-b}\ A\ B\ C = A + t *_{\mathbb{R}} (C - A)$
unfolding *conway-A-b-def*
using *conway-point-past-first-on-line[of A C dist B C]* **by** *blast*
show
 $\exists t::\text{real. } \text{conway-C-b}\ A\ B\ C = A + t *_{\mathbb{R}} (C - A)$
unfolding *conway-C-b-def*
using *conway-point-past-second-on-line[of A C dist A B]* **by** *blast*
show
 $\exists t::\text{real. } \text{conway-B-a}\ A\ B\ C = B + t *_{\mathbb{R}} (C - B)$
unfolding *conway-B-a-def*
using *conway-point-past-first-on-line[of B C dist C A]* **by** *blast*

```

show
   $\exists t :: \text{real}. \text{conway-C-a } A \ B \ C = B + t *_{\mathbb{R}} (C - B)$ 
unfolding conway-C-a-def
using conway-point-past-second-on-line[of B C dist A B] by blast
show
   $\text{dist } A \ (\text{conway-A-c } A \ B \ C) = \text{dist } B \ C$ 
   $\text{dist } B \ (\text{conway-B-c } A \ B \ C) = \text{dist } C \ A$ 
   $\text{dist } A \ (\text{conway-A-b } A \ B \ C) = \text{dist } B \ C$ 
   $\text{dist } C \ (\text{conway-C-b } A \ B \ C) = \text{dist } A \ B$ 
   $\text{dist } B \ (\text{conway-B-a } A \ B \ C) = \text{dist } C \ A$ 
   $\text{dist } C \ (\text{conway-C-a } A \ B \ C) = \text{dist } A \ B$ 
using conway-points-have-declared-lengths[OF sides] by blast+
qed

definition conway-points ::  $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow (\text{real}^2) \text{ set}$ 
where conway-points  $A \ B \ C =$ 
   $\{\text{conway-A-c } A \ B \ C, \text{conway-B-c } A \ B \ C, \text{conway-A-b } A \ B \ C,$ 
   $\text{conway-C-b } A \ B \ C, \text{conway-B-a } A \ B \ C, \text{conway-C-a } A \ B \ C\}$ 

corollary conway-circle-theorem-sphere:
fixes  $A \ B \ C :: \text{real}^2$ 
assumes triangle:  $\neg \text{collinear } \{A, B, C\}$ 
shows
   $\text{conway-points } A \ B \ C$ 
   $\subseteq \text{sphere } (\text{triangle-incenter } A \ B \ C)$ 
   $\quad (\text{sqrt } ((\text{triangle-inradius } A \ B \ C)^2$ 
   $\quad + (\text{triangle-semiperimeter } A \ B \ C)^2))$ 
using conway-circle-theorem[OF triangle]
unfolding conway-points-def sphere-def by auto

corollary conway-circle-theorem-set:
fixes  $A \ B \ C :: \text{real}^2$ 
assumes triangle:  $\neg \text{collinear } \{A, B, C\}$ 
shows
   $\forall P \in \text{conway-points } A \ B \ C.$ 
   $\text{dist } (\text{triangle-incenter } A \ B \ C) \ P$ 
   $= \text{sqrt } ((\text{triangle-inradius } A \ B \ C)^2$ 
   $\quad + (\text{triangle-semiperimeter } A \ B \ C)^2)$ 
using conway-circle-theorem[OF triangle]
unfolding conway-points-def by auto

end

```

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