

Interval Temporal Logic on Natural Numbers

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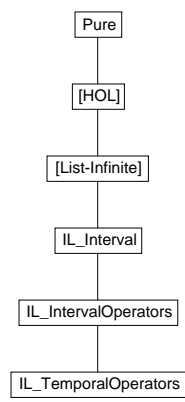
Abstract

We introduce a theory of temporal logic operators using sets of natural numbers as time domain, formalized in a shallow embedding manner. The theory comprises special natural intervals (theory `IL.Interval`: open and closed intervals, continuous and modulo intervals, interval traversing results), operators for shifting intervals to left/right on the number axis as well as expanding/contracting intervals by constant factors (theory `IL.IntervalOperators.thy`), and ultimately definitions and results for unary and binary temporal operators on arbitrary natural sets (theory `IL.TemporalOperators`).

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1 IL-Interval: Intervals and operations for temporal logic declarations

```

theory IL-Interval
imports
  Main Nat-Infinity
  SetInterval2 InfiniteSet2 SetIntervalStep
  Util-Nat Util-MinMax Util-Div
begin

```

1.1 Time intervals – definitions and basic lemmata

1.1.1 Definitions

```

types Time = nat

```

```

types iT = Time set
typ iT

```

Infinite interval starting at some natural n .

```

definition
  iFROM :: Time  $\Rightarrow$  iT ([...]
where
  [...]  $\equiv$  {n..}

```

Finite interval starting at $0::'a$ and ending at some natural n .

```

definition
  iTILL :: Time  $\Rightarrow$  iT ([...]
where
  [...n]  $\equiv$  {...n}

```

Finite bounded interval containing the naturals between n and $n + d$. d denotes the difference between left and right interval bound. The number of elements is $d + (1::'a)$ so that an empty interval cannot be defined.

```

definition
  iIN :: Time  $\Rightarrow$  nat  $\Rightarrow$  iT ([...,])
where
  [n..,d]  $\equiv$  {n..n+d}

```

Infinite modulo interval containing all naturals having the same division remainder modulo m as r , and beginning at n .

```

definition
  iMOD :: Time  $\Rightarrow$  nat  $\Rightarrow$  iT ([-, mod -])
where
  [r, mod m]  $\equiv$  { x. x mod m = r mod m  $\wedge$  r  $\leq$  x }

```

Finite bounded modulo interval containing all naturals having the same division remainder modulo m as r , beginning at n , and ending after c cycles

at $r + m * c$. The number of elements is $c + (1::'a)$ so that an empty interval cannot be defined.

definition

$iMODb :: Time \Rightarrow nat \Rightarrow nat \Rightarrow iT ([-, mod -, -])$

where

$[r, mod m, c] \equiv \{ x. x \text{ mod } m = r \text{ mod } m \wedge r \leq x \wedge x \leq r + m * c \}$

1.1.2 Membership in an interval

lemmas $iT-defs = iFROM-def \ iTILL-def \ iIN-def \ iMOD-def \ iMODb-def$

lemma $iFROM-iff: x \in [n..] = (n \leq x)$

$\langle proof \rangle$

lemma $iTILL-iff: x \in [..n] = (x \leq n)$

$\langle proof \rangle$

lemma $iIN-iff: x \in [n..,d] = (n \leq x \wedge x \leq n + d)$

$\langle proof \rangle$

lemma $iMOD-iff: x \in [r, mod m] = (x \text{ mod } m = r \text{ mod } m \wedge r \leq x)$

$\langle proof \rangle$

lemma $iMODb-iff: x \in [r, mod m, c] =$

$(x \text{ mod } m = r \text{ mod } m \wedge r \leq x \wedge x \leq r + m * c)$

$\langle proof \rangle$

lemma $iFROM-D: x \in [n..] \Longrightarrow (n \leq x)$

$\langle proof \rangle$

lemma $iTILL-D: x \in [..n] \Longrightarrow (x \leq n)$

$\langle proof \rangle$

corollary $iIN-geD: x \in [n..,d] \Longrightarrow n \leq x$

$\langle proof \rangle$

corollary $iIN-leD: x \in [n..,d] \Longrightarrow x \leq n + d$

$\langle proof \rangle$

corollary $iMOD-modD: x \in [r, mod m] \Longrightarrow x \text{ mod } m = r \text{ mod } m$

$\langle proof \rangle$

corollary $iMOD-geD: x \in [r, mod m] \Longrightarrow r \leq x$

$\langle proof \rangle$

corollary $iMODb-modD: x \in [r, mod m, c] \Longrightarrow x \text{ mod } m = r \text{ mod } m$

$\langle proof \rangle$

corollary $iMODb-geD: x \in [r, mod m, c] \Longrightarrow r \leq x$

$\langle proof \rangle$

corollary $iMODb-leD: x \in [r, mod m, c] \Longrightarrow x \leq r + m * c$

$\langle proof \rangle$

lemmas $iT-iff = iFROM-iff \ iTILL-iff \ iIN-iff \ iMOD-iff \ iMODb-iff$

lemmas $iT-drule =$

$iFROM-D$

$iTILL-D$

$iIN-geD \ iIN-leD$

$iMOD-modD \ iMOD-geD$

iMODb-modD iMODb-geD iMODb-leD

thm *iT-drule*

lemma

iFROM-I [intro]: $n \leq x \implies x \in [n..]$ **and**

iTILL-I [intro]: $x \leq n \implies x \in [..n]$ **and**

iIN-I [intro]: $n \leq x \implies x \leq n + d \implies x \in [n..,d]$ **and**

iMOD-I [intro]: $x \bmod m = r \bmod m \implies r \leq x \implies x \in [r, \bmod m]$ **and**

iMODb-I [intro]: $x \bmod m = r \bmod m \implies r \leq x \implies x \leq r + m * c \implies x \in [r, \bmod m, c]$
 ⟨proof⟩

lemma

iFROM-E [elim]: $x \in [n..] \implies (n \leq x \implies P) \implies P$ **and**

iTILL-E [elim]: $x \in [..n] \implies (x \leq n \implies P) \implies P$ **and**

iIN-E [elim]: $x \in [n..,d] \implies (n \leq x \implies x \leq n + d \implies P) \implies P$ **and**

iMOD-E [elim]: $x \in [r, \bmod m] \implies (x \bmod m = r \bmod m \implies r \leq x \implies P) \implies P$ **and**

iMODb-E [elim]: $x \in [r, \bmod m, c] \implies (x \bmod m = r \bmod m \implies r \leq x \implies x \leq r + m * c \implies P) \implies P$
 ⟨proof⟩

lemma *iIN-Suc-insert-conv*:

insert (*Suc* ($n + d$)) [$n..,d$] = [$n..,Suc\ d$]

⟨proof⟩

lemma *iTILL-Suc-insert-conv*: *insert* (*Suc* n) [$..n$] = [$..Suc\ n$]

⟨proof⟩

lemma *iMODb-Suc-insert-conv*:

insert ($r + m * Suc\ c$) [$r, \bmod m, c$] = [$r, \bmod m, Suc\ c$]

⟨proof⟩

thm *less-mod-eq-imp-add-divisor-le*

⟨proof⟩

thm *add-le-imp-le-right*

⟨proof⟩

lemma *iFROM-pred-insert-conv*: *insert* ($n - Suc\ 0$) [$n..$] = [$n - Suc\ 0..$]

⟨proof⟩

lemma *iIN-pred-insert-conv*:

$0 < n \implies insert\ (n - Suc\ 0)\ [n..,d] = [n - Suc\ 0..,Suc\ d]$

⟨proof⟩

lemma *iMOD-pred-insert-conv*:

$m \leq r \implies \text{insert } (r - m) [r, \text{mod } m] = [r - m, \text{mod } m]$
 ⟨proof⟩

thm *less-mod-eq-imp-add-divisor-le*[of $r - m$ x m]
 ⟨proof⟩

lemma *iMODb-pred-insert-conv*:

$m \leq r \implies \text{insert } (r - m) [r, \text{mod } m, c] = [r - m, \text{mod } m, \text{Suc } c]$
 ⟨proof⟩

lemma *iFROM-Suc-pred-insert-conv*: $\text{insert } n [\text{Suc } n \dots] = [n \dots]$
 ⟨proof⟩

lemma *iIN-Suc-pred-insert-conv*: $\text{insert } n [\text{Suc } n \dots, d] = [n \dots, \text{Suc } d]$
 ⟨proof⟩

lemma *iMOD-Suc-pred-insert-conv*: $\text{insert } r [r + m, \text{mod } m] = [r, \text{mod } m]$
 ⟨proof⟩

lemma *iMODb-Suc-pred-insert-conv*: $\text{insert } r [r + m, \text{mod } m, c] = [r, \text{mod } m, \text{Suc } c]$
 ⟨proof⟩

lemmas *iT-Suc-insert* =
iIN-Suc-insert-conv
iTILL-Suc-insert-conv
iMODb-Suc-insert-conv

lemmas *iT-pred-insert* =
iFROM-pred-insert-conv
iIN-pred-insert-conv
iMOD-pred-insert-conv
iMODb-pred-insert-conv

lemmas *iT-Suc-pred-insert* =
iFROM-Suc-pred-insert-conv
iIN-Suc-pred-insert-conv
iMOD-Suc-pred-insert-conv
iMODb-Suc-pred-insert-conv

lemma *iMOD-mem-diff*: $\llbracket a \in [r, \text{mod } m]; b \in [r, \text{mod } m] \rrbracket \implies (a - b) \text{ mod } m = 0$

⟨proof⟩

lemma *iMODb-mem-diff*: $\llbracket a \in [r, \text{mod } m, c]; b \in [r, \text{mod } m, c] \rrbracket \implies (a - b) \text{ mod } m = 0$

⟨proof⟩

1.1.3 Interval conversions

lemma *iIN-0-iTILL-conv*: $[0 \dots, n] = [\dots n]$
 ⟨proof⟩

lemma *iIN-iTILL-iTILL-conv*: $0 < n \implies [n \dots, d] = [\dots n + d] - [\dots n - \text{Suc } 0]$
 ⟨proof⟩

lemma *iIN-iFROM-iTILL-conv*: $[n \dots, d] = [n \dots] \cap [\dots n + d]$
 ⟨proof⟩

lemma *iMODb-iMOD-iTILL-conv*: $[r, \text{mod } m, c] = [r, \text{mod } m] \cap [\dots r + m * c]$
 ⟨proof⟩

lemma *iMODb-iMOD-iIN-conv*: $[r, \text{mod } m, c] = [r, \text{mod } m] \cap [r \dots, m * c]$
 ⟨proof⟩

lemma *iFROM-iTILL-iIN-conv*: $n \leq n' \implies [n \dots] \cap [\dots n'] = [n \dots, n' - n]$
 ⟨proof⟩

lemma *iMOD-iTILL-iMODb-conv*:

$r \leq n \implies [r, \text{mod } m] \cap [\dots n] = [r, \text{mod } m, (n - r) \text{ div } m]$

⟨proof⟩

thm *le-imp-sub-mod-le*

⟨proof⟩

lemma *iMOD-iIN-iMODb-conv*:

$[r, \text{mod } m] \cap [r \dots, d] = [r, \text{mod } m, d \text{ div } m]$

⟨proof⟩

thm *iMOD-iTILL-iMODb-conv*

⟨proof⟩

thm *subst[of {}] - λt. ∀x.(x - t) = x*

thm *subst[of {}] - λt. ∀x.(x - t) = x, THEN spec*

⟨proof⟩

thm *disjoint-iff-not-equal*

⟨proof⟩

thm *UNIV-def*

lemma *iFROM-0*: $[0 \dots] = \text{UNIV}$

⟨proof⟩

lemma *iTILL-0*: $[\dots 0] = \{0\}$

⟨proof⟩

lemma *iIN-0*: $[n \dots, 0] = \{n\}$

⟨proof⟩

lemma *iMOD-0*: $[r, \text{mod } 0] = [r \dots, 0]$

⟨proof⟩

lemma *iMODb-mod-0*: $[r, \text{mod } 0, c] = [r \dots, 0]$

⟨proof⟩

lemma *iMODb-0*: $[r, \text{mod } m, 0] = [r \dots, 0]$

⟨proof⟩

lemmas *iT-0 =*

iFROM-0

iTILL-0

iIN-0

iMOD-0

iMODb-mod-0

iMODb-0

thm *iT-0*

lemma *iMOD-1*: $[r, \text{mod } \text{Suc } 0] = [r \dots]$

<proof>

lemma *iMODb-mod-1*: $[r, \text{mod } \text{Suc } 0, c] = [r..c]$

<proof>

1.1.4 Finiteness and emptiness of intervals

lemma

iFROM-not-empty: $[n..] \neq \{\}$ **and**

iTILL-not-empty: $[..n] \neq \{\}$ **and**

iIN-not-empty: $[n..d] \neq \{\}$ **and**

iMOD-not-empty: $[r, \text{mod } m] \neq \{\}$ **and**

iMODb-not-empty: $[r, \text{mod } m, c] \neq \{\}$

<proof>

lemmas *iT-not-empty* =

iFROM-not-empty

iTILL-not-empty

iIN-not-empty

iMOD-not-empty

iMODb-not-empty

thm *iT-not-empty*

lemma

iTILL-finite: *finite* $[..n]$ **and**

iIN-finite: *finite* $[n..d]$ **and**

iMODb-finite: *finite* $[r, \text{mod } m, c]$ **and**

iMOD-0-finite: *finite* $[r, \text{mod } 0]$

<proof>

lemma *iFROM-infinite*: *infinite* $[n..]$

<proof>

lemma *iMOD-infinite*: $0 < m \implies \text{infinite } [r, \text{mod } m]$

thm *infinite-nat-iff-asc-chain*

<proof>

lemmas *iT-finite* =

iTILL-finite

iIN-finite

iMODb-finite *iMOD-0-finite*

thm *iT-finite*

lemmas *iT-infinite* =

iFROM-infinite

iMOD-infinite

thm *iT-infinite*

thm

iMax-finite-conv
iMax-infinite-conv

1.1.5 Min and Max element of an interval**lemma**

iTILL-Min: $iMin [\dots n] = 0$ **and**
iFROM-Min: $iMin [n \dots] = n$ **and**
iIN-Min: $iMin [n \dots, d] = n$ **and**
iMOD-Min: $iMin [r, \text{mod } m] = r$ **and**
iMODb-Min: $iMin [r, \text{mod } m, c] = r$

thm *iMin-equality**\langle proof \rangle***lemmas** *iT-Min =*

iIN-Min
iTILL-Min
iFROM-Min
iMOD-Min
iMODb-Min

thm *iT-Min***lemma**

iTILL-Max: $Max [\dots n] = n$ **and**
iIN-Max: $Max [n \dots, d] = n+d$ **and**
iMODb-Max: $Max [r, \text{mod } m, c] = r + m * c$ **and**
iMOD-0-Max: $Max [r, \text{mod } 0] = r$

*\langle proof \rangle***lemmas** *iT-Max =*

iTILL-Max
iIN-Max
iMODb-Max
iMOD-0-Max

thm *iT-Max***lemma**

iTILL-iMax: $iMax [\dots n] = Fin\ n$ **and**
iIN-iMax: $iMax [n \dots, d] = Fin\ (n+d)$ **and**
iMODb-iMax: $iMax [r, \text{mod } m, c] = Fin\ (r + m * c)$ **and**
iMOD-0-iMax: $iMax [r, \text{mod } 0] = Fin\ r$ **and**
iFROM-iMax: $iMax [n \dots] = \infty$ **and**
iMOD-iMax: $0 < m \implies iMax [r, \text{mod } m] = \infty$

*\langle proof \rangle***lemmas** *iT-iMax =*

iTILL-iMax
iIN-iMax

iMODb-iMax
iMOD-0-iMax
iFROM-iMax
iMOD-iMax
thm *iT-iMax*

1.2 Adding and subtracting constants to interval elements

lemma

iFROM-plus: $x \in [n\dots] \implies x + k \in [n\dots]$ **and**
iFROM-Suc: $x \in [n\dots] \implies \text{Suc } x \in [n\dots]$ **and**
iFROM-minus: $\llbracket x \in [n\dots]; k \leq x - n \rrbracket \implies x - k \in [n\dots]$ **and**
iFROM-pred: $n < x \implies x - \text{Suc } 0 \in [n\dots]$
 ⟨proof⟩

lemma

iTILL-plus: $\llbracket x \in [\dots n]; k \leq n - x \rrbracket \implies x + k \in [\dots n]$ **and**
iTILL-Suc: $x < n \implies \text{Suc } x \in [\dots n]$ **and**
iTILL-minus: $x \in [\dots n] \implies x - k \in [\dots n]$ **and**
iTILL-pred: $x \in [\dots n] \implies x - \text{Suc } 0 \in [\dots n]$
 ⟨proof⟩

lemma *iIN-plus*: $\llbracket x \in [n\dots, d]; k \leq n + d - x \rrbracket \implies x + k \in [n\dots, d]$
 ⟨proof⟩

lemma *iIN-Suc*: $\llbracket x \in [n\dots, d]; x < n + d \rrbracket \implies \text{Suc } x \in [n\dots, d]$
 ⟨proof⟩

lemma *iIN-minus*: $\llbracket x \in [n\dots, d]; k \leq x - n \rrbracket \implies x - k \in [n\dots, d]$
 ⟨proof⟩

lemma *iIN-pred*: $\llbracket x \in [n\dots, d]; n < x \rrbracket \implies x - \text{Suc } 0 \in [n\dots, d]$
 ⟨proof⟩

lemma *iMOD-plus-divisor-mult*: $x \in [r, \text{mod } m] \implies x + k * m \in [r, \text{mod } m]$
 ⟨proof⟩

corollary *iMOD-plus-divisor*: $x \in [r, \text{mod } m] \implies x + m \in [r, \text{mod } m]$
 ⟨proof⟩

lemma *iMOD-minus-divisor-mult*:

$\llbracket x \in [r, \text{mod } m]; k * m \leq x - r \rrbracket \implies x - k * m \in [r, \text{mod } m]$

thm *mod-diff-mult-self1*

⟨proof⟩

corollary *iMOD-minus-divisor-mult2*:

$\llbracket x \in [r, \text{mod } m]; k \leq (x - r) \text{ div } m \rrbracket \implies x - k * m \in [r, \text{mod } m]$

⟨proof⟩

thm *mod-0-div-mult-cancel*[THEN *iffD1*, OF *mod-eq-imp-diff-mod-0*]

⟨proof⟩

corollary *iMOD-minus-divisor*:

$\llbracket x \in [r, \text{mod } m]; m + r \leq x \rrbracket \implies x - m \in [r, \text{mod } m]$

⟨proof⟩

thm *iMOD-minus-divisor-mult*[of *x r m 1*]

⟨proof⟩

lemma *iMOD-plus*:

$$x \in [r, \text{mod } m] \implies (x + k \in [r, \text{mod } m]) = (k \text{ mod } m = 0)$$

<proof>

thm *mod-add-eq-imp-mod-0*[THEN *iffD1*]

<proof>

corollary *iMOD-Suc*:

$$x \in [r, \text{mod } m] \implies (\text{Suc } x \in [r, \text{mod } m]) = (m = \text{Suc } 0)$$

<proof>

lemma *iMOD-minus*:

$$\llbracket x \in [r, \text{mod } m]; k \leq x - r \rrbracket \implies (x - k \in [r, \text{mod } m]) = (k \text{ mod } m = 0)$$

<proof>

corollary *iMOD-pred*:

$$\llbracket x \in [r, \text{mod } m]; r < x \rrbracket \implies (x - \text{Suc } 0 \in [r, \text{mod } m]) = (m = \text{Suc } 0)$$

<proof>

thm *iMOD-Suc*[of $x - \text{Suc } 0$, THEN *iffD1*]

<proof>

lemma *iMODb-plus-divisor-mult*:

$$\llbracket x \in [r, \text{mod } m, c]; k * m \leq r + m * c - x \rrbracket \implies x + k * m \in [r, \text{mod } m, c]$$

<proof>

lemma *iMODb-plus-divisor-mult2*:

$$\llbracket x \in [r, \text{mod } m, c]; k \leq c - (x - r) \text{ div } m \rrbracket \implies$$

$$x + k * m \in [r, \text{mod } m, c]$$

<proof>

lemma *iMODb-plus-divisor*:

$$\llbracket x \in [r, \text{mod } m, c]; x < r + m * c \rrbracket \implies x + m \in [r, \text{mod } m, c]$$

thm *less-mod-eq-imp-add-divisor-le*

<proof>

lemma *iMODb-minus-divisor-mult*:

$$\llbracket x \in [r, \text{mod } m, c]; r + k * m \leq x \rrbracket \implies x - k * m \in [r, \text{mod } m, c]$$

thm *mod-diff-mult-self1*

<proof>

lemma *iMODb-plus*:

$$\llbracket x \in [r, \text{mod } m, c]; k \leq r + m * c - x \rrbracket \implies$$

$$(x + k \in [r, \text{mod } m, c]) = (k \text{ mod } m = 0)$$

<proof>

thm *mod-add-eq-imp-mod-0*[THEN *iffD1*]

<proof>

corollary *iMODb-Suc*:

$$\llbracket x \in [r, \text{mod } m, c]; x < r + m * c \rrbracket \implies$$

$$(\text{Suc } x \in [r, \text{mod } m, c]) = (m = \text{Suc } 0)$$

<proof>

lemma *iMODb-minus*:

$$\llbracket x \in [r, \text{mod } m, c]; k \leq x - r \rrbracket \implies$$

$$(x - k \in [r, \text{mod } m, c]) = (k \text{ mod } m = 0)$$

<proof>

corollary *iMODb-pred*:

$$\llbracket x \in [r, \text{mod } m, c]; r < x \rrbracket \implies$$

$(x - \text{Suc } 0 \in [r, \text{mod } m, c]) = (m = \text{Suc } 0)$
 ⟨proof⟩
thm *iMOD-pred*[*THEN iffD1, of x r m*]
 ⟨proof⟩

lemmas *iFROM-plus-minus* =
iFROM-plus
iFROM-Suc
iFROM-minus
iFROM-pred
thm *iFROM-plus-minus*

lemmas *iTILL-plus-minus* =
iTILL-plus
iTILL-Suc
iTILL-minus
iTILL-pred
thm *iTILL-plus-minus*

lemmas *iIN-plus-minus* =
iIN-plus
iIN-Suc
iTILL-minus
iIN-pred
thm *iIN-plus-minus*

lemmas *iMOD-plus-minus-divisor* =
iMOD-plus-divisor-mult
iMOD-plus-divisor
iMOD-minus-divisor-mult
iMOD-minus-divisor-mult2
iMOD-minus-divisor
thm *iMOD-plus-minus-divisor*

lemmas *iMOD-plus-minus* =
iMOD-plus
iMOD-Suc
iMOD-minus
iMOD-pred
thm *iMOD-plus-minus*

lemmas *iMODb-plus-minus-divisor* =
iMODb-plus-divisor-mult
iMODb-plus-divisor-mult2
iMODb-plus-divisor
iMODb-minus-divisor-mult
thm *iMODb-plus-minus-divisor*

lemmas *iMODb-plus-minus* =

iMODb-plus
iMODb-Suc
iMODb-minus
iMODb-pred

thm *iMODb-plus-minus*

lemmas *iT-plus-minus* =
iFROM-plus-minus
iTILL-plus-minus
iIN-plus-minus
iMOD-plus-minus-divisor
iMOD-plus-minus
iMODb-plus-minus-divisor
iMODb-plus-minus

thm *iT-plus-minus*

1.3 Relations between intervals

1.3.1 Auxiliary lemmata

lemma *Suc-in-imp-not-subset-iMOD*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies \neg S \subseteq [r, \text{mod } m]$

thm *iMOD-Suc[THEN iffD1]*

<proof>

corollary *Suc-in-imp-neq-iMOD*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies S \neq [r, \text{mod } m]$

<proof>

lemma *Suc-in-imp-not-subset-iMODb*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies \neg S \subseteq [r, \text{mod } m, c]$

<proof>

thm *iMODb-Suc[THEN iffD1]*

<proof>

corollary *Suc-in-imp-neq-iMODb*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies S \neq [r, \text{mod } m, c]$

<proof>

1.3.2 Subset relation between intervals

lemma

iIN-iFROM-subset-same: $[n\dots,d] \subseteq [n\dots]$ **and**

iIN-iTILL-subset-same: $[n\dots,d] \subseteq [\dots,n+d]$ **and**

iMOD-iFROM-subset-same: $[r, \text{mod } m] \subseteq [r\dots]$ **and**

iMODb-iTILL-subset-same: $[r, \text{mod } m, c] \subseteq [\dots,r+m*c]$ **and**

iMODb-iIN-subset-same: $[r, \text{mod } m, c] \subseteq [r\dots,m*c]$ **and**

iMODb-iMOD-subset-same: $[r, \text{mod } m, c] \subseteq [r, \text{mod } m]$

<proof>

lemmas *iT-subset-same* =

iIN-iFROM-subset-same
iIN-iTILL-subset-same
iMOD-iFROM-subset-same
iMODb-iTILL-subset-same
iMODb-iIN-subset-same
iMODb-iTILL-subset-same
iMODb-iMOD-subset-same
thm *iT-subset-same*

lemma *iMODb-imp-iMOD*: $x \in [r, \text{mod } m, c] \implies x \in [r, \text{mod } m]$
 ⟨proof⟩

lemma *iMOD-imp-iMODb*:
 $\llbracket x \in [r, \text{mod } m]; x \leq r + m * c \rrbracket \implies x \in [r, \text{mod } m, c]$
 ⟨proof⟩

lemma *iMOD-singleton-subset-conv*: $([r, \text{mod } m] \subseteq \{a\}) = (r = a \wedge m = 0)$
 ⟨proof⟩

lemma *iMOD-singleton-eq-conv*: $([r, \text{mod } m] = \{a\}) = (r = a \wedge m = 0)$
 ⟨proof⟩

lemma *iMODb-singleton-subset-conv*:
 $([r, \text{mod } m, c] \subseteq \{a\}) = (r = a \wedge (m = 0 \vee c = 0))$
 ⟨proof⟩

lemma *iMODb-singleton-eq-conv*:
 $([r, \text{mod } m, c] = \{a\}) = (r = a \wedge (m = 0 \vee c = 0))$
 ⟨proof⟩

lemma *iMODb-subset-imp-divisor-mod-0*:
 $\llbracket 0 < c'; [r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c] \rrbracket \implies m' \text{ mod } m = 0$
 ⟨proof⟩

thm *mod-add-eq-imp-mod-0*[of $r' m' m$]
 ⟨proof⟩

lemma *iMOD-subset-imp-divisor-mod-0*:
 $[r', \text{mod } m'] \subseteq [r, \text{mod } m] \implies m' \text{ mod } m = 0$
 ⟨proof⟩

thm *mod-add-eq-imp-mod-0*[of $r' m' m$]
 ⟨proof⟩

lemma *iMOD-subset-imp-iMODb-subset*:
 $\llbracket [r', \text{mod } m'] \subseteq [r, \text{mod } m]; r' + m' * c' \leq r + m * c \rrbracket \implies$
 $[r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]$
 ⟨proof⟩

lemma *iMODb-subset-imp-iMOD-subset*:

$$\llbracket [r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]; 0 < c' \rrbracket \implies \\ [r', \text{mod } m'] \subseteq [r, \text{mod } m]$$

thm *subsetD*

<proof>

thm *subsetI*

<proof>

thm *mod-eq-mod-0-imp-mod-eq*

<proof>

thm *iMODb-subset-imp-divisor-mod-0*

<proof>

lemma *iMODb-0-iMOD-subset-conv*:

$$([r', \text{mod } m', 0] \subseteq [r, \text{mod } m]) = \\ (r' \text{ mod } m = r \text{ mod } m \wedge r \leq r')$$

<proof>

lemma *iFROM-subset-conv*: $([n'..] \subseteq [n..]) = (n \leq n')$

<proof>

lemma *iFROM-iMOD-subset-conv*: $([n'..] \subseteq [r, \text{mod } m]) = (r \leq n' \wedge m = \text{Suc } 0)$

<proof>

thm *iMin-subset[OF iFROM-not-empty]*

<proof>

thm *Suc-in-imp-not-subset-iMOD*

<proof>

lemma *iIN-subset-conv*: $([n'..,d'] \subseteq [n..,d]) = (n \leq n' \wedge n'+d' \leq n+d)$

<proof>

lemma *iIN-iFROM-subset-conv*: $([n'..,d'] \subseteq [n..]) = (n \leq n')$

<proof>

lemma *iIN-iTILL-subset-conv*: $([n'..,d'] \subseteq [..n]) = (n' + d' \leq n)$

<proof>

lemma *iIN-iMOD-subset-conv*:

$$0 < d' \implies ([n'..,d'] \subseteq [r, \text{mod } m]) = (r \leq n' \wedge m = \text{Suc } 0)$$

<proof>

thm *Suc-in-imp-not-subset-iMOD*

<proof>

lemma *iIN-iMODb-subset-conv*:

$$0 < d' \implies \\ ([n'..,d'] \subseteq [r, \text{mod } m, c]) = \\ (r \leq n' \wedge m = \text{Suc } 0 \wedge n' + d' \leq r + m * c)$$

<proof>

thm *subset-trans[OF - iMODb-iMOD-subset-same]*

$\langle proof \rangle$

lemma *iTILL-subset-conv*: $([..n^\frown] \subseteq [..n]) = (n' \leq n)$

$\langle proof \rangle$

lemma *iTILL-iFROM-subset-conv*: $([..n^\frown] \subseteq [n..]) = (n = 0)$

$\langle proof \rangle$

lemma *iTILL-iIN-subset-conv*: $([..n^\frown] \subseteq [n\dots,d]) = (n = 0 \wedge n' \leq d)$

$\langle proof \rangle$

lemma *iTILL-iMOD-subset-conv*:

$0 < n' \implies ([..n^\frown] \subseteq [r, \text{mod } m]) = (r = 0 \wedge m = \text{Suc } 0)$

$\langle proof \rangle$

lemma *iTILL-iMODb-subset-conv*:

$0 < n' \implies ([..n^\frown] \subseteq [r, \text{mod } m, c]) = (r = 0 \wedge m = \text{Suc } 0 \wedge n' \leq r + m * c)$

$\langle proof \rangle$

lemma *iMOD-iFROM-subset-conv*: $([r', \text{mod } m^\frown] \subseteq [n\dots]) = (n \leq r')$

$\langle proof \rangle$

lemma *iMODb-iFROM-subset-conv*: $([r', \text{mod } m', c^\frown] \subseteq [n\dots]) = (n \leq r')$

$\langle proof \rangle$

lemma *iMODb-iIN-subset-conv*:

$([r', \text{mod } m', c^\frown] \subseteq [n\dots,d]) = (n \leq r' \wedge r' + m' * c' \leq n + d)$

$\langle proof \rangle$

lemma *iMODb-iTILL-subset-conv*:

$([r', \text{mod } m', c^\frown] \subseteq [..n]) = (r' + m' * c' \leq n)$

$\langle proof \rangle$

lemma *iMOD-0-subset-conv*: $([r', \text{mod } 0] \subseteq [r, \text{mod } m]) = (r' \text{ mod } m = r \text{ mod } m \wedge r \leq r')$

$\langle proof \rangle$

lemma *iMOD-subset-conv*: $0 < m \implies$

$([r', \text{mod } m^\frown] \subseteq [r, \text{mod } m]) =$

$(r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0)$

$\langle proof \rangle$

thm *mod-eq-mod-0-imp-mod-eq*

$\langle proof \rangle$

lemma *iMODb-subset-mod-0-conv*:

$([r', \text{mod } m', c^\frown] \subseteq [r, \text{mod } 0, c]) = (r'=r \wedge (m'=0 \vee c'=0))$

$\langle proof \rangle$

lemma *iMODb-subset-0-conv*:

$([r', \text{mod } m', c^\frown] \subseteq [r, \text{mod } m, 0]) = (r'=r \wedge (m'=0 \vee c'=0))$

$\langle proof \rangle$

lemma *iMODb-0-subset-conv*:

$([r', \text{mod } m', 0] \subseteq [r, \text{mod } m, c]) = (r' \in [r, \text{mod } m, c])$

$\langle proof \rangle$

lemma *iMODb-mod-0-subset-conv*:

$([r', \text{mod } 0, c'] \subseteq [r, \text{mod } m, c]) = (r' \in [r, \text{mod } m, c])$
 <proof>

lemma *iMODb-subset-conv'*: $\llbracket 0 < c; 0 < c' \rrbracket \implies$
 $([r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]) =$
 $(r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0 \wedge$
 $r' + m' * c' \leq r + m * c)$
 <proof>

thm *iMODb-subset-imp-iMOD-subset*
 <proof>

thm *iMOD-subset-imp-iMODb-subset*
 <proof>

lemma *iMODb-subset-conv*: $\llbracket 0 < m'; 0 < c' \rrbracket \implies$
 $([r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]) =$
 $(r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0 \wedge$
 $r' + m' * c' \leq r + m * c)$
 <proof>

thm *add-less-le-mono*[of 0 m' * c' r r']
 <proof>

lemma *iMODb-iMOD-subset-conv*: $0 < c' \implies$
 $([r', \text{mod } m', c'] \subseteq [r, \text{mod } m]) =$
 $(r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0)$
 <proof>

thm *subsetD*[OF - iMinI-ex2[OF iMODb-not-empty]]
 <proof>

thm *iMOD-iTILL-iMODb-conv* *iMODb-subset-imp-divisor-mod-0*
 <proof>

thm *subset-trans*[OF iMODb-iMOD-subset-same]
 <proof>

lemmas *iT-subset-conv* =
iFROM-subset-conv
iFROM-iMOD-subset-conv
iTILL-subset-conv
iTILL-iFROM-subset-conv
iTILL-iIN-subset-conv
iTILL-iMOD-subset-conv
iTILL-iMODb-subset-conv
iIN-subset-conv
iIN-iFROM-subset-conv
iIN-iTILL-subset-conv
iIN-iMOD-subset-conv
iIN-iMODb-subset-conv
iMOD-subset-conv
iMOD-iFROM-subset-conv

iMODb-subset-conv'
iMODb-subset-conv
iMODb-iFROM-subset-conv
iMODb-iIN-subset-conv
iMODb-iTILL-subset-conv
iMODb-iMOD-subset-conv

thm *iT-subset-conv*

lemma *iFROM-subset*: $n \leq n' \implies [n'..] \subseteq [n..]$

<proof>

lemma *not-iFROM-iIN-subset*: $\neg [n'..] \subseteq [n..,d]$

<proof>

lemma *not-iFROM-iTILL-subset*: $\neg [n'..] \subseteq [..n]$

<proof>

lemma *not-iFROM-iMOD-subset*: $m \neq \text{Suc } 0 \implies \neg [n'..] \subseteq [r, \text{mod } m]$

<proof>

lemma *not-iFROM-iMODb-subset*: $\neg [n'..] \subseteq [r, \text{mod } m, c]$

thm *infinite-not-subset-finite*

<proof>

lemma *iIN-subset*: $\llbracket n \leq n'; n' + d' \leq n + d \rrbracket \implies [n'..,d'] \subseteq [n..,d]$

<proof>

lemma *iIN-iFROM-subset*: $n \leq n' \implies [n'..,d'] \subseteq [n..]$

<proof>

lemma *iIN-iTILL-subset*: $n' + d' \leq n \implies [n'..,d'] \subseteq [..n]$

<proof>

lemma *not-iIN-iMODb-subset*: $\llbracket 0 < d'; m \neq \text{Suc } 0 \rrbracket \implies \neg [n'..,d'] \subseteq [r, \text{mod } m, c]$

<proof>

lemma *not-iIN-iMOD-subset*: $\llbracket 0 < d'; m \neq \text{Suc } 0 \rrbracket \implies \neg [n'..,d'] \subseteq [r, \text{mod } m]$

<proof>

thm *iIN-iTILL-subset[OF order-refl]*

thm *Int-greatest[OF - iIN-iTILL-subset[OF order-refl]]*

<proof>

thm *iMOD-iTILL-iMODb-conv not-iIN-iMODb-subset*

<proof>

lemma *iTILL-subset*: $n' \leq n \implies [..n'] \subseteq [..n]$

<proof>

lemma *iTILL-iFROM-subset*: $([..n'] \subseteq [0..])$

<proof>

lemma *iTILL-iIN-subset*: $n' \leq d \implies ([..n'] \subseteq [0..,d])$

<proof>

thm *not-iIN-iMOD-subset*

lemma *not-iTILL-iMOD-subset*:

$\llbracket 0 < n'; m \neq \text{Suc } 0 \rrbracket \implies \neg [..n'] \subseteq [r, \text{mod } m]$

<proof>

lemma *not-iTILL-iMODb-subset*:

$\llbracket 0 < n'; m \neq \text{Suc } 0 \rrbracket \implies \neg [\dots n'] \subseteq [r, \text{mod } m, c]$
 <proof>

lemma *iMOD-iFROM-subset*: $n \leq r' \implies [r', \text{mod } m'] \subseteq [n\dots]$
 <proof>

lemma *not-iMOD-iIN-subset*: $0 < m' \implies \neg [r', \text{mod } m'] \subseteq [n\dots, d]$
 <proof>

lemma *not-iMOD-iTILL-subset*: $0 < m' \implies \neg [r', \text{mod } m'] \subseteq [\dots n]$
 <proof>

thm *iMOD-subset-conv*

lemma *iMOD-subset*:

$\llbracket r \leq r'; r' \text{ mod } m = r \text{ mod } m; m' \text{ mod } m = 0 \rrbracket \implies [r', \text{mod } m'] \subseteq [r, \text{mod } m]$
 <proof>

lemma *not-iMOD-iMODb-subset*: $0 < m' \implies \neg [r', \text{mod } m'] \subseteq [r, \text{mod } m, c]$
 <proof>

lemma *iMODb-iFROM-subset*: $n \leq r' \implies [r', \text{mod } m', c'] \subseteq [n\dots]$

thm *iMODb-iFROM-subset-conv* [THEN iffD2]
 <proof>

lemma *iMODb-iTILL-subset*:

$r' + m' * c' \leq n \implies [r', \text{mod } m', c'] \subseteq [\dots n]$
 <proof>

thm *iMODb-iIN-subset-conv*

lemma *iMODb-iIN-subset*:

$\llbracket n \leq r'; r' + m' * c' \leq n + d \rrbracket \implies [r', \text{mod } m', c'] \subseteq [n\dots, d]$
 <proof>

thm *iMODb-iMOD-subset-conv*

lemma *iMODb-iMOD-subset*:

$\llbracket r \leq r'; r' \text{ mod } m = r \text{ mod } m; m' \text{ mod } m = 0 \rrbracket \implies [r', \text{mod } m', c'] \subseteq [r, \text{mod } m]$

<proof>

thm *iMODb-iMOD-subset-conv*

<proof>

thm *iMODb-subset-conv*

lemma *iMODb-subset*:

$\llbracket r \leq r'; r' \text{ mod } m = r \text{ mod } m; m' \text{ mod } m = 0; r' + m' * c' \leq r + m * c \rrbracket \implies$
 $[r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]$
 <proof>

lemma *iFROM-trans*: $\llbracket y \in [x\dots]; z \in [y\dots] \rrbracket \implies z \in [x\dots]$

<proof>

lemma *iTILL-trans*: $\llbracket y \in [\dots x]; z \in [\dots y] \rrbracket \implies z \in [\dots x]$

<proof>

thm *iIN-subset*

lemma *iIN-trans*:

$$\llbracket y \in [x \dots, d]; z \in [y \dots, d']; d' \leq x + d - y \rrbracket \implies z \in [x \dots, d]$$

<proof>

lemma *iMOD-trans*:

$$\llbracket y \in [x, \text{mod } m]; z \in [y, \text{mod } m] \rrbracket \implies z \in [x, \text{mod } m]$$

<proof>

lemma *iMODb-trans*:

$$\llbracket y \in [x, \text{mod } m, c]; z \in [y, \text{mod } m, c']; m * c' \leq x + m * c - y \rrbracket \implies z \in [x, \text{mod } m, c]$$

<proof>

lemma *iMODb-trans'*:

$$\llbracket y \in [x, \text{mod } m, c]; z \in [y, \text{mod } m, c']; c' \leq x \text{ div } m + c - y \text{ div } m \rrbracket \implies z \in [x, \text{mod } m, c]$$

<proof>

1.3.3 Equality of intervals

lemma *iFROM-eq-conv*: $([n \dots] = [n' \dots]) = (n = n')$

<proof>

lemma *iIN-eq-conv*: $([n \dots, d] = [n' \dots, d']) = (n = n' \wedge d = d')$

<proof>

lemma *iTILL-eq-conv*: $([\dots n] = [\dots n']) = (n = n')$

thm *iIN-eq-conv[of 0 n 0 n']*

<proof>

thm *iMOD-singleton-eq-conv*

lemma *iMOD-0-eq-conv*: $([r, \text{mod } 0] = [r', \text{mod } m']) = (r = r' \wedge m' = 0)$

<proof>

thm *iMOD-singleton-eq-conv*

<proof>

lemma *iMOD-eq-conv*: $0 < m \implies ([r, \text{mod } m] = [r', \text{mod } m']) = (r = r' \wedge m = m')$

<proof>

thm *iMODb-singleton-eq-conv*

lemma *iMODb-mod-0-eq-conv*:

$$([r, \text{mod } 0, c] = [r', \text{mod } m', c']) = (r = r' \wedge (m' = 0 \vee c' = 0))$$

<proof>

lemma *iMODb-0-eq-conv*:

$$([r, \text{mod } m, 0] = [r', \text{mod } m', c']) = (r = r' \wedge (m' = 0 \vee c' = 0))$$

<proof>

lemma *iMODb-eq-conv*: $\llbracket 0 < m; 0 < c \rrbracket \implies$

$$([r, \text{mod } m, c] = [r', \text{mod } m', c']) = (r = r' \wedge m = m' \wedge c = c')$$

<proof>

lemma *iMOD-iFROM-eq-conv*: $([n\dots] = [r, \text{mod } m]) = (n = r \wedge m = \text{Suc } 0)$

<proof>

thm *iMODb-singleton-eq-conv*

lemma *iMODb-iIN-0-eq-conv*:

$([n\dots, 0] = [r, \text{mod } m, c]) = (n = r \wedge (m = 0 \vee c = 0))$

<proof>

lemma *iMODb-iIN-eq-conv*:

$0 < d \implies ([n\dots, d] = [r, \text{mod } m, c]) = (n = r \wedge m = \text{Suc } 0 \wedge c = d)$

<proof>

1.3.4 Inequality of intervals

lemma *iFROM-iIN-neq*: $[n'\dots] \neq [n\dots, d]$

<proof>

corollary *iFROM-iTILL-neq*: $[n'\dots] \neq [\dots n]$

<proof>

corollary *iFROM-iMOD-neq*: $m \neq \text{Suc } 0 \implies [n\dots] \neq [r, \text{mod } m]$

<proof>

corollary *iFROM-iMODb-neq*: $[n\dots] \neq [r, \text{mod } m, c]$

<proof>

corollary *iIN-iMOD-neq*: $0 < m \implies [n\dots, d] \neq [r, \text{mod } m]$

<proof>

corollary *iIN-iMODb-neq2*: $\llbracket m \neq \text{Suc } 0; 0 < d \rrbracket \implies [n\dots, d] \neq [r, \text{mod } m, c]$

<proof>

lemma *iIN-iMODb-neq*: $\llbracket 2 \leq m; 0 < c \rrbracket \implies [n\dots, d] \neq [r, \text{mod } m, c]$

<proof>

thm *iMODb-singleton-eq-conv*

<proof>

lemma *iTILL-iIN-neq*: $0 < n \implies [\dots n] \neq [n\dots, d]$

<proof>

corollary *iTILL-iMOD-neq*: $0 < m \implies [\dots n] \neq [r, \text{mod } m]$

<proof>

corollary *iTILL-iMODb-neq*:

$\llbracket m \neq \text{Suc } 0; 0 < n \rrbracket \implies [\dots n] \neq [r, \text{mod } m, c]$

<proof>

lemma *iMOD-iMODb-neq*: $0 < m \implies [r, \text{mod } m] \neq [r', \text{mod } m', c]$

<proof>

lemmas *iT-neq* =

iFROM-iTILL-neq *iFROM-iIN-neq* *iFROM-iMOD-neq* *iFROM-iMODb-neq*

iTILL-iIN-neq *iTILL-iMOD-neq* *iTILL-iMODb-neq*

iIN-iMOD-neq *iIN-iMODb-neq* *iIN-iMODb-neq2*

iMOD-iMODb-neq

thm *iT-neq*

1.4 Union and intersection of intervals

lemma *iFROM-union'*: $[n\dots] \cup [n'\dots] = [\min n n'\dots]$

<proof>

corollary *iFROM-union*: $n \leq n' \implies [n\dots] \cup [n'\dots] = [n\dots]$

<proof>

lemma *iFROM-inter'*: $[n\dots] \cap [n'\dots] = [\max n n'\dots]$

<proof>

corollary *iFROM-inter*: $n' \leq n \implies [n\dots] \cap [n'\dots] = [n\dots]$

<proof>

lemma *iTILL-union'*: $[\dots n] \cup [\dots n'] = [\dots \max n n']$

<proof>

corollary *iTILL-union*: $n' \leq n \implies [\dots n] \cup [\dots n'] = [\dots n]$

<proof>

lemma *iTILL-iFROM-union*: $n \leq n' \implies [\dots n'] \cup [n\dots] = UNIV$

<proof>

lemma *iTILL-inter'*: $[\dots n] \cap [\dots n'] = [\dots \min n n']$

<proof>

corollary *iTILL-inter*: $n \leq n' \implies [\dots n] \cap [\dots n'] = [\dots n]$

<proof>

Union and intersection for iIN only when there intersection is not empty and none of them is other's subset, for instance: .. n .. n+d .. n' .. n'+d'

lemma *iIN-union*:

$[[n \leq n'; n' \leq \text{Suc}(n + d); n + d \leq n' + d']] \implies$

$[n\dots, d] \cup [n'\dots, d'] = [n\dots, n' - n + d']$

<proof>

lemma *iIN-append-union*:

$[n\dots, d] \cup [n + d\dots, d'] = [n\dots, d + d']$

<proof>

lemma *iIN-append-union-Suc*:

$[n\dots, d] \cup [\text{Suc}(n + d)\dots, d'] = [n\dots, \text{Suc}(d + d')]$

<proof>

lemma *iIN-append-union-pred*:

$0 < d \implies [n\dots, d - \text{Suc } 0] \cup [n + d\dots, d'] = [n\dots, d + d']$

<proof>

lemma *iIN-iFROM-union*:

$n' \leq \text{Suc}(n + d) \implies [n\dots, d] \cup [n'\dots] = [\min n n'\dots]$

<proof>

lemma *iIN-iFROM-append-union*:

$[n\dots, d] \cup [n + d\dots] = [n\dots]$

<proof>

lemma *iIN-iFROM-append-union-Suc*:

$$[n\dots,d] \cup [\text{Suc } (n + d)\dots] = [n\dots]$$

<proof>

lemma *iIN-iFROM-append-union-pred*:

$$0 < d \implies [n\dots,d - \text{Suc } 0] \cup [n + d\dots] = [n\dots]$$

<proof>

lemma *iIN-inter*:

$$\begin{aligned} & \llbracket n \leq n'; n' \leq n + d; n + d \leq n' + d' \rrbracket \implies \\ & [n\dots,d] \cap [n'\dots,d'] = [n'\dots,n + d - n'] \end{aligned}$$

<proof>

lemma *iMOD-union*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m \rrbracket \implies \\ & [r, \bmod m] \cup [r', \bmod m] = [r, \bmod m] \end{aligned}$$

<proof>

lemma *iMOD-union'*:

$$\begin{aligned} & r \bmod m = r' \bmod m \implies \\ & [r, \bmod m] \cup [r', \bmod m] = [\min r r', \bmod m] \end{aligned}$$

<proof>

lemma *iMOD-inter*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m \rrbracket \implies \\ & [r, \bmod m] \cap [r', \bmod m] = [r', \bmod m] \end{aligned}$$

<proof>

lemma *iMOD-inter'*:

$$\begin{aligned} & r \bmod m = r' \bmod m \implies \\ & [r, \bmod m] \cap [r', \bmod m] = [\max r r', \bmod m] \end{aligned}$$

<proof>

lemma *iMODb-union*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m; r' \leq r + m * c; r + m * c \leq r' + m * c' \rrbracket \implies \\ & [r, \bmod m, c] \cup [r', \bmod m, c'] = [r, \bmod m, r' \text{ div } m - r \text{ div } m + c'] \end{aligned}$$

<proof>

thm *iMODb-iMOD-subset-same*

thm *Un-absorb1[OF iMODb-iMOD-subset-same]*

lemma *iMODb-append-union*:

$$[r, \bmod m, c] \cup [r + m * c, \bmod m, c'] = [r, \bmod m, c + c']$$

thm *iMODb-union[of r r + m * c m c c']*

<proof>

lemma *iMODb-iMOD-append-union'*:

$$\begin{aligned} & \llbracket r \bmod m = r' \bmod m; r' \leq r + m * \text{Suc } c \rrbracket \implies \\ & [r, \bmod m, c] \cup [r', \bmod m] = [\min r r', \bmod m] \end{aligned}$$

<proof>

thm *add-le-imp-le-right*[of - m]

<proof>

thm *less-mod-eq-imp-add-divisor-le*

<proof>

lemma *iMODb-iMOD-append-union*:

$$\llbracket r \leq r'; r \bmod m = r' \bmod m; r' \leq r + m * \text{Suc } c \rrbracket \implies$$

$$[r, \bmod m, c] \cup [r', \bmod m] = [r, \bmod m]$$

thm *iMODb-iMOD-append-union'*

<proof>

lemma *iMODb-append-union-Suc*:

$$[r, \bmod m, c] \cup [r + m * \text{Suc } c, \bmod m, c'] = [r, \bmod m, \text{Suc } (c + c')]$$

thm *insert-absorb*[of $r + m * c$ $[r, \bmod m, c] \cup [r + m * \text{Suc } c, \bmod m, c']$]

<proof>

thm *iMODb-append-union*[of r m c]

<proof>

lemma *iMODb-append-union-pred*:

$$0 < c \implies [r, \bmod m, c - \text{Suc } 0] \cup [r + m * c, \bmod m, c'] = [r, \bmod m, c + c']$$

<proof>

lemma *iMODb-inter*:

$$\llbracket r \leq r'; r \bmod m = r' \bmod m; r' \leq r + m * c; r + m * c \leq r' + m * c' \rrbracket \implies$$

$$[r, \bmod m, c] \cap [r', \bmod m, c'] = [r', \bmod m, c - (r' - r) \text{ div } m]$$

<proof>

thm *mod-0-div-mult-cancel*[THEN *iffD1*, OF *mod-eq-imp-diff-mod-0*]

<proof>

lemmas *iT-union'* =

iFROM-union'

iTILL-union'

iMOD-union'

iMODb-iMOD-append-union'

lemmas *iT-union* =

iFROM-union

iTILL-union

iTILL-iFROM-union

iIN-union

iIN-iFROM-union

iMOD-union

iMODb-union

lemmas *iT-union-append* =

iIN-append-union

iIN-append-union-Suc

iIN-append-union-pred
iIN-iFROM-append-union
iIN-iFROM-append-union-Suc
iIN-iFROM-append-union-pred
iMODb-append-union
iMODb-iMOD-append-union
iMODb-append-union-Suc
iMODb-append-union-pred

lemmas *iT-inter'* =
iFROM-inter'
iTILL-inter'
iMOD-inter'

lemmas *iT-inter* =
iFROM-inter
iTILL-inter
iIN-inter
iMOD-inter
iMODb-inter

thm *iT-union'*
thm *iT-union*
thm *iT-union-append*
thm *iT-inter'*
thm *iT-inter*

thm *partition-Union*

lemma *mod-partition-Union*:

$$0 < m \implies (\bigcup k. A \cap [k * m \dots, m - \text{Suc } 0]) = A$$

<proof>

thm *subst[where s=UNIV and P= $\lambda x. A \cap x = A$]*

<proof>

thm *div-mult-cancel*

<proof>

thm *le-add-diff*

<proof>

thm *Suc-mod-le-divisor*

<proof>

thm *mod-partition-Union*

lemma *finite-mod-partition-Union*:

$$\llbracket 0 < m; \text{finite } A \rrbracket \implies$$

$$(\bigcup k \leq \text{Max } A \text{ div } m. A \cap [k * m \dots, m - \text{Suc } 0]) = A$$

thm *subst[OF mod-partition-Union[of m], where*

$$P = \lambda x. (\bigcup k \leq \text{Max } A \text{ div } m. A \cap [k * m \dots, m - \text{Suc } 0]) = x]$$

<proof>

thm *setsum-UN-disjoint*

lemma *mod-partition-is-disjoint*:

$\llbracket 0 < (m::nat); k \neq k' \rrbracket \implies$

$(A \cap [k * m.., m - \text{Suc } 0]) \cap (A \cap [k' * m.., m - \text{Suc } 0]) = \{\}$

<proof>

thm *le-less-imp-div*

<proof>

1.5 Cutting intervals

thm *i-cut-defs*

thm *i-cut-subset*

thm *i-cut-bound*

thm

cut-le-less-conv

cut-less-le-conv

thm

cut-ge-greater-conv

cut-greater-ge-conv

thm

cut-le-Min-empty

cut-less-Min-empty

cut-le-Min-not-empty

cut-less-Min-not-empty

thm

i-cut-min-empty

thm

i-cut-min-all

thm

i-cut-max-empty

thm

i-cut-max-all

lemma *iTILL-cut-le*: $[\dots n] \downarrow \leq t = (\text{if } t \leq n \text{ then } [\dots t] \text{ else } [\dots n])$

<proof>

corollary *iTILL-cut-le1*: $t \in [\dots n] \implies [\dots n] \downarrow \leq t = [\dots t]$

<proof>

corollary *iTILL-cut-le2*: $t \notin [\dots n] \implies [\dots n] \downarrow \leq t = [\dots n]$

<proof>

lemma *iFROM-cut-le*:

$[\dots] \downarrow \leq t =$

$(\text{if } t < n \text{ then } \{\} \text{ else } [n.., t-n])$

<proof>

corollary *iFROM-cut-le1*: $t \in [n\dots] \implies [n\dots] \downarrow \leq t = [n\dots, t - n]$

<proof>

lemma *iIN-cut-le*:

$[n\dots, d] \downarrow \leq t =$
if $t < n$ then $\{\}$ else
if $t \leq n + d$ then $[n\dots, t - n]$
else $[n\dots, d]$

<proof>

corollary *iIN-cut-le1*:

$t \in [n\dots, d] \implies [n\dots, d] \downarrow \leq t = [n\dots, t - n]$

<proof>

lemma *iMOD-cut-le*:

$[r, \text{mod } m] \downarrow \leq t =$
if $t < r$ then $\{\}$
else $[r, \text{mod } m, (t - r) \text{ div } m]$

<proof>

thm *less-mod-eq-imp-add-divisor-le*

<proof>

thm *mod-eq-imp-diff-mod-eq*[of - m r t, rule-format]

<proof>

thm *add-le-mono2*[OF mod-le-divisor]

thm *order-trans*[OF add-le-mono2[OF mod-le-divisor]]

<proof>

thm *le-imp-sub-mod-le*[of - t]

<proof>

lemma *iMOD-cut-le1*:

$t \in [r, \text{mod } m] \implies$
 $[r, \text{mod } m] \downarrow \leq t = [r, \text{mod } m, (t - r) \text{ div } m]$

<proof>

lemma *iMODb-cut-le*:

$[r, \text{mod } m, c] \downarrow \leq t =$
if $t < r$ then $\{\}$
else if $t < r + m * c$ then $[r, \text{mod } m, (t - r) \text{ div } m]$
else $[r, \text{mod } m, c]$

<proof>

thm *iMOD-iTILL-iMODb-conv*[of r r + m * c m]

<proof>

find-theorems *cut-le* (- \cap -)

<proof>

thm *iMOD-iTILL-iMODb-conv*

\langle proof \rangle

lemma *iMODb-cut-le1*:

$t \in [r, \text{mod } m, c] \implies$
 $[r, \text{mod } m, c] \downarrow \leq t = [r, \text{mod } m, (t - r) \text{ div } m]$
 \langle proof \rangle

lemma *iTILL-cut-less*:

$[..n] \downarrow < t = ($
 $\text{if } n < t \text{ then } [..n] \text{ else}$
 $\text{if } t = 0 \text{ then } \{\}$
 $\text{else } [..t - \text{Suc } 0])$
 \langle proof \rangle

lemma *iTILL-cut-less1*:

$\llbracket t \in [..n]; 0 < t \rrbracket \implies [..n] \downarrow < t = [..t - \text{Suc } 0]$
thm *iTILL-cut-less*
 \langle proof \rangle

lemma *iFROM-cut-less*:

$[n..] \downarrow < t = ($
 $\text{if } t \leq n \text{ then } \{\}$
 $\text{else } [n.., t - \text{Suc } n])$
 \langle proof \rangle

lemma *iFROM-cut-less1*:

$n < t \implies [n..] \downarrow < t = [n.., t - \text{Suc } n]$
 \langle proof \rangle

lemma *iIN-cut-less*:

$[n.., d] \downarrow < t = ($
 $\text{if } t \leq n \text{ then } \{\} \text{ else}$
 $\text{if } t \leq n + d \text{ then } [n.., t - \text{Suc } n]$
 $\text{else } [n.., d])$
 \langle proof \rangle

lemma *iIN-cut-less1*:

$\llbracket t \in [n.., d]; n < t \rrbracket \implies [n.., d] \downarrow < t = [n.., t - \text{Suc } n]$
thm *iIN-cut-less*
 \langle proof \rangle

thm

cut-le-less-conv
cut-less-le-conv

lemma *iMOD-cut-less*:

$$[r, \text{mod } m] \downarrow < t = ($$

$$\text{if } t \leq r \text{ then } \{\}$$

$$\text{else } [r, \text{mod } m, (t - \text{Suc } r) \text{ div } m])$$

thm *iMOD-cut-le*

<proof>

lemma *iMOD-cut-less1*:

$$\llbracket t \in [r, \text{mod } m]; r < t \rrbracket \implies$$

$$[r, \text{mod } m] \downarrow < t = [r, \text{mod } m, (t - r) \text{ div } m - \text{Suc } 0]$$

<proof>

thm *mod-0-imp-diff-Suc-div-conv mod-eq-imp-diff-mod-0*

<proof>

lemma *iMODb-cut-less*:

$$[r, \text{mod } m, c] \downarrow < t = ($$

$$\text{if } t \leq r \text{ then } \{\} \text{ else}$$

$$\text{if } r + m * c < t \text{ then } [r, \text{mod } m, c]$$

$$\text{else } [r, \text{mod } m, (t - \text{Suc } r) \text{ div } m])$$

thm *iMODb-cut-le*

<proof>

lemma *iMODb-cut-less1*: $\llbracket t \in [r, \text{mod } m, c]; r < t \rrbracket \implies$

$$[r, \text{mod } m, c] \downarrow < t = [r, \text{mod } m, (t - r) \text{ div } m - \text{Suc } 0]$$

<proof>

thm *mod-0-imp-diff-Suc-div-conv mod-eq-imp-diff-mod-0*

<proof>

lemmas *iT-cut-le* =

iTILL-cut-le
iFROM-cut-le
iIN-cut-le
iMOD-cut-le
iMODb-cut-le

thm *iT-cut-le*

lemmas *iT-cut-le1* =

iTILL-cut-le1
iFROM-cut-le1
iIN-cut-le1
iMOD-cut-le1
iMODb-cut-le1

thm *iT-cut-le1*

lemmas *iT-cut-less* =

iTILL-cut-less
iFROM-cut-less
iIN-cut-less

iMOD-cut-less
iMODb-cut-less
thm *iT-cut-less*
lemmas *iT-cut-less1* =
iTILL-cut-less1
iFROM-cut-less1
iIN-cut-less1
iMOD-cut-less1
iMODb-cut-less1
thm *iT-cut-less1*

lemmas *iT-cut-le-less* =
iTILL-cut-le
iTILL-cut-less
iFROM-cut-le
iFROM-cut-less
iIN-cut-le
iIN-cut-less
iMOD-cut-le
iMOD-cut-less
iMODb-cut-le
iMODb-cut-less
lemmas *iT-cut-le-less1* =
iTILL-cut-le1
iTILL-cut-less1
iFROM-cut-le1
iFROM-cut-less1
iIN-cut-le1
iIN-cut-less1
iMOD-cut-le1
iMOD-cut-less1
iMODb-cut-le1
iMODb-cut-less1
thm *iT-cut-le-less*
thm *iT-cut-le-less1*

lemma *iTILL-cut-ge*:
 $[\dots n] \downarrow \geq t = (\text{if } n < t \text{ then } \{\} \text{ else } [t\dots, n-t])$
<proof>
corollary *iTILL-cut-ge1*: $t \in [\dots n] \implies [\dots n] \downarrow \geq t = [t\dots, n-t]$
<proof>
corollary *iTILL-cut-ge2*: $t \notin [\dots n] \implies [\dots n] \downarrow \geq t = \{\}$
<proof>

lemma *iTILL-cut-greater*:
 $[\dots n] \downarrow > t = (\text{if } n \leq t \text{ then } \{\} \text{ else } [Suc\ t\dots, n - Suc\ t])$
<proof>
corollary *iTILL-cut-greater1*:
 $t \in [\dots n] \implies t < n \implies [\dots n] \downarrow > t = [Suc\ t\dots, n - Suc\ t]$

<proof>

corollary *iTILL-cut-greater2*: $t \notin [\dots n] \implies [\dots n] \downarrow > t = \{\}$

<proof>

lemma *iFROM-cut-ge*:

$[\dots] \downarrow \geq t = (\text{if } n \leq t \text{ then } [t\dots] \text{ else } [n\dots])$

<proof>

corollary *iFROM-cut-ge1*: $t \in [n\dots] \implies [n\dots] \downarrow \geq t = [t\dots]$

<proof>

lemma *iFROM-cut-greater*:

$[\dots] \downarrow > t = (\text{if } n \leq t \text{ then } [\text{Suc } t\dots] \text{ else } [n\dots])$

<proof>

corollary *iFROM-cut-greater1*:

$t \in [n\dots] \implies [n\dots] \downarrow > t = [\text{Suc } t\dots]$

<proof>

lemma *iIN-cut-ge*:

$[\dots, d] \downarrow \geq t = (\text{if } t < n \text{ then } [n\dots, d] \text{ else } \text{if } t \leq n+d \text{ then } [t\dots, n+d-t] \text{ else } \{\})$

<proof>

corollary *iIN-cut-ge1*: $t \in [n\dots, d] \implies$

$[\dots, d] \downarrow \geq t = [t\dots, n+d-t]$

<proof>

corollary *iIN-cut-ge2*: $t \notin [n\dots, d] \implies$

$[\dots, d] \downarrow \geq t = (\text{if } t < n \text{ then } [n\dots, d] \text{ else } \{\})$

<proof>

lemma *iIN-cut-greater*:

$[\dots, d] \downarrow > t = (\text{if } t < n \text{ then } [n\dots, d] \text{ else } \text{if } t < n+d \text{ then } [\text{Suc } t\dots, n+d-\text{Suc } t] \text{ else } \{\})$

<proof>

corollary *iIN-cut-greater1*:

$\llbracket t \in [n\dots, d]; t < n+d \rrbracket \implies$

$[\dots, d] \downarrow > t = [\text{Suc } t\dots, n+d-\text{Suc } t]$

<proof>

thm *sub-diff-mod-eq*^[of r t 1 m, simplified]

lemma *mod-cut-greater-aux-t-less*:

$$\llbracket 0 < (m::nat); r \leq t \rrbracket \implies \\ t < t + m - (t - r) \bmod m$$

thm *less-add-diff*

<proof>

lemma *mod-cut-greater-aux-le-x*:

$$\llbracket (r::nat) \leq t; t < x; x \bmod m = r \bmod m \rrbracket \implies \\ t + m - (t - r) \bmod m \leq x$$

thm *diff-mod-le*

<proof>

thm *diff-add-assoc2*^[of (t - r) mod m t m]

<proof>

thm *less-mod-eq-imp-add-divisor-le*^[of t - (t - r) mod m x m]

<proof>

thm *sub-diff-mod-eq*

<proof>

lemma *iMOD-cut-greater*:

$$[r, \bmod m] \downarrow > t = (\\ \text{if } t < r \text{ then } [r, \bmod m] \text{ else} \\ \text{if } m = 0 \text{ then } \{\} \\ \text{else } [t + m - (t - r) \bmod m, \bmod m])$$

<proof>

thm *order-trans*^[OF mod-le-dividend]

<proof>

thm *diff-add-assoc2*^[of (t - r) mod m t m]

<proof>

thm *sub-diff-mod-eq*

<proof>

thm *less-mod-eq-imp-add-divisor-le*^[of t - (t - r) mod m x m]

<proof>

lemma *iMOD-cut-greater1*:

$$t \in [r, \bmod m] \implies \\ [r, \bmod m] \downarrow > t = (\\ \text{if } m = 0 \text{ then } \{\} \\ \text{else } [t + m, \bmod m])$$

<proof>

lemma *iMODb-cut-greater-aux*:

$$\llbracket 0 < m; t < r + m * c; r \leq t \rrbracket \implies \\ (r + m * c - (t + m - (t - r) \bmod m)) \operatorname{div} m = \\ c - \operatorname{Suc} ((t - r) \operatorname{div} m)$$

thm *diff-diff-right*^[of r t+m-(t-r)mod m m*c, symmetric]

<proof>

thm *mod-cut-greater-aux-t-less*

<proof>

lemma *iMODb-cut-greater*:

$$[r, \text{mod } m, c] \downarrow > t = ($$

if $t < r$ then $[r, \text{mod } m, c]$ else

if $r + m * c \leq t$ then $\{\}$

else $[t + m - (t - r) \text{ mod } m, \text{mod } m, c - \text{Suc } ((t-r) \text{ div } m)]$)

<proof>

thm *iMOD-iTILL-iMODb-conv*[of $r \ r + m * c \ m$]

<proof>

thm *iMOD-cut-greater*

<proof>

thm *iMOD-iTILL-iMODb-conv*

<proof>

thm *mod-cut-greater-aux-le-x*

<proof>

lemma *iMODb-cut-greater1*:

$$t \in [r, \text{mod } m, c] \implies$$

$$[r, \text{mod } m, c] \downarrow > t = ($$

if $r + m * c \leq t$ then $\{\}$

else $[t + m, \text{mod } m, c - \text{Suc } ((t-r) \text{ div } m)]$)

<proof>

lemma *iMOD-cut-ge*:

$$[r, \text{mod } m] \downarrow \geq t = ($$

if $t \leq r$ then $[r, \text{mod } m]$ else

if $m = 0$ then $\{\}$

else $[t + m - \text{Suc } ((t - \text{Suc } r) \text{ mod } m), \text{mod } m]$)

<proof>

thm *iMOD-cut-greater nat-cut-greater-ge-conv*[*symmetric*]

<proof>

lemma *iMOD-cut-ge1*:

$$t \in [r, \text{mod } m] \implies$$

$$[r, \text{mod } m] \downarrow \geq t = [t, \text{mod } m]$$

<proof>

lemma *iMODb-cut-ge*:

$$[r, \text{mod } m, c] \downarrow \geq t = ($$

if $t \leq r$ then $[r, \text{mod } m, c]$ else

if $r + m * c < t$ then $\{\}$

else $[t + m - \text{Suc } ((t - \text{Suc } r) \text{ mod } m), \text{mod } m, c - (t + m - \text{Suc } r) \text{ div } m]$)

thm *iMOD-cut-ge*

<proof>

thm *diff-mod-pred*

<proof>

thm *mod-0-imp-diff-Suc-div-conv*
 ⟨proof⟩
thm *div-le-mono*
 ⟨proof⟩
thm *diff-mod-pred*
 ⟨proof⟩
thm *iMOD-iTILL-iMODb-conv*[of r $r + m * c$ m]
 ⟨proof⟩
thm *iMOD-cut-ge*
 ⟨proof⟩
thm *iMOD-iTILL-iMODb-conv*
 ⟨proof⟩
thm *mod-cut-greater-aux-le-x*
 ⟨proof⟩
thm *iMODb-cut-greater-aux*[of m t r c]
 ⟨proof⟩
thm *div-diff1-eq*
 ⟨proof⟩

thm *iMODb-cut-greater1*
lemma *iMODb-cut-ge1*:
 $t \in [r, \text{mod } m, c] \implies$
 $[r, \text{mod } m, c] \downarrow \geq t =$
 $\text{if } r + m * c < t \text{ then } \{\}$
 $\text{else } [t, \text{mod } m, c - (t - r) \text{ div } m]$

⟨proof⟩
thm *iMODb-cut-ge*
 ⟨proof⟩
thm *mod-eq-imp-diff-mod-eq-divisor*
 ⟨proof⟩
thm *div-add1-eq*
 ⟨proof⟩

lemma *iMOD-0-cut-greater*:
 $t \in [r, \text{mod } 0] \implies [r, \text{mod } 0] \downarrow > t = \{\}$
 ⟨proof⟩

lemma *iMODb-0-cut-greater*: $t \in [r, \text{mod } 0, c] \implies$
 $[r, \text{mod } 0, c] \downarrow > t = \{\}$
 ⟨proof⟩

lemmas *iT-cut-ge* =
iTILL-cut-ge
iFROM-cut-ge
iIN-cut-ge
iMOD-cut-ge
iMODb-cut-ge

thm *iT-cut-ge*
lemmas *iT-cut-ge1* =
iTILL-cut-ge1
iFROM-cut-ge1

iIN-cut-ge1
iMOD-cut-ge1
iMODb-cut-ge1
thm *iT-cut-le1*

lemmas *iT-cut-greater* =
iTILL-cut-greater
iFROM-cut-greater
iIN-cut-greater
iMOD-cut-greater
iMODb-cut-greater

thm *iT-cut-greater*

lemmas *iT-cut-greater1* =
iTILL-cut-greater1
iFROM-cut-greater1
iIN-cut-greater1
iMOD-cut-greater1
iMODb-cut-greater1

thm *iT-cut-greater1*

lemmas *iT-cut-ge-greater* =
iTILL-cut-ge
iTILL-cut-greater
iFROM-cut-ge
iFROM-cut-greater
iIN-cut-ge
iIN-cut-greater
iMOD-cut-ge
iMOD-cut-greater
iMODb-cut-ge
iMODb-cut-greater

lemmas *iT-cut-ge-greater1* =
iTILL-cut-ge1
iTILL-cut-greater1
iFROM-cut-ge1
iFROM-cut-greater1
iIN-cut-ge1
iIN-cut-greater1
iMOD-cut-ge1
iMOD-cut-greater1
iMODb-cut-ge1
iMODb-cut-greater1

thm *iT-cut-ge-greater*

thm *iT-cut-ge-greater1*

1.6 Cardinality of intervals

lemma *iFROM-card*: $\text{card } [n..] = 0$
<proof>

lemma *iTILL-card*: $\text{card } [\dots n] = \text{Suc } n$
 ⟨proof⟩
lemma *iIN-card*: $\text{card } [n\dots, d] = \text{Suc } d$
 ⟨proof⟩
lemma *iMOD-0-card*: $\text{card } [r, \text{mod } 0] = \text{Suc } 0$
 ⟨proof⟩
lemma *iMOD-card*: $0 < m \implies \text{card } [r, \text{mod } m] = 0$
 ⟨proof⟩
lemma *iMOD-card-if*: $\text{card } [r, \text{mod } m] = (\text{if } m = 0 \text{ then } \text{Suc } 0 \text{ else } 0)$
 ⟨proof⟩
lemma *iMODb-mod-0-card*: $\text{card } [r, \text{mod } 0, c] = \text{Suc } 0$
 ⟨proof⟩
lemma *iMODb-card*: $0 < m \implies \text{card } [r, \text{mod } m, c] = \text{Suc } c$
 ⟨proof⟩
thm *iMODb-Suc-insert-conv*
 ⟨proof⟩
lemma *iMODb-card-if*:
 $\text{card } [r, \text{mod } m, c] = (\text{if } m = 0 \text{ then } \text{Suc } 0 \text{ else } \text{Suc } c)$
 ⟨proof⟩

lemmas *iT-card* =
iFROM-card
iTILL-card
iIN-card
iMOD-card-if
iMODb-card-if

Cardinality with *icard*

lemma *iFROM-icard*: $\text{icard } [n\dots] = \infty$
 ⟨proof⟩
lemma *iTILL-icard*: $\text{icard } [\dots n] = \text{Fin } (\text{Suc } n)$
 ⟨proof⟩
lemma *iIN-icard*: $\text{icard } [n\dots, d] = \text{Fin } (\text{Suc } d)$
 ⟨proof⟩
lemma *iMOD-0-icard*: $\text{icard } [r, \text{mod } 0] = \text{iSuc } 0$
 ⟨proof⟩
lemma *iMOD-icard*: $0 < m \implies \text{icard } [r, \text{mod } m] = \infty$
 ⟨proof⟩
lemma *iMOD-icard-if*: $\text{icard } [r, \text{mod } m] = (\text{if } m = 0 \text{ then } \text{iSuc } 0 \text{ else } \infty)$
 ⟨proof⟩
lemma *iMODb-mod-0-icard*: $\text{icard } [r, \text{mod } 0, c] = \text{iSuc } 0$
 ⟨proof⟩
lemma *iMODb-icard*: $0 < m \implies \text{icard } [r, \text{mod } m, c] = \text{Fin } (\text{Suc } c)$
 ⟨proof⟩
lemma *iMODb-icard-if*: $\text{icard } [r, \text{mod } m, c] = \text{Fin } (\text{if } m = 0 \text{ then } \text{Suc } 0 \text{ else } \text{Suc } c)$
 ⟨proof⟩

lemmas *iT-icard* =

iFROM-icard
iTILL-icard
iIN-icard
iMOD-icard-if
iMODb-icard-if

1.7 Functions *inext* and *iprev* with intervals

thm

inext-def
iprev-def

thm

inext-Max
iprev-iMin

thm

inext-closed
iprev-closed

thm

iMin-subset
iMin-Un

thm

Max-Un
Min-Un

lemma

iFROM-inext: $t \in [n..] \implies \text{inext } t [n..] = \text{Suc } t$ **and**
iTILL-inext: $t < n \implies \text{inext } t [..n] = \text{Suc } t$ **and**
iIN-inext: $\llbracket n \leq t; t < n + d \rrbracket \implies \text{inext } t [n..,d] = \text{Suc } t$
 <proof>

lemma

iFROM-iprev': $t \in [n..] \implies \text{iprev } (\text{Suc } t) [n..] = t$ **and**
iFROM-iprev: $n < t \implies \text{iprev } t [n..] = t - \text{Suc } 0$ **and**
iTILL-iprev: $t \in [..n] \implies \text{iprev } t [..n] = t - \text{Suc } 0$ **and**
iIN-iprev: $\llbracket n < t; t \leq n + d \rrbracket \implies \text{iprev } t [n..,d] = t - \text{Suc } 0$ **and**
iIN-iprev': $\llbracket n \leq t; t < n + d \rrbracket \implies \text{iprev } (\text{Suc } t) [n..,d] = t$
 <proof>

lemma *iMOD-inext*: $t \in [r, \text{mod } m] \implies \text{inext } t [r, \text{mod } m] = t + m$
 <proof>

lemma *iMOD-iprev*: $\llbracket t \in [r, \text{mod } m]; r < t \rrbracket \implies \text{iprev } t [r, \text{mod } m] = t - m$
 <proof>

thm *mod-eq-imp-diff-mod-eq-divisor*

<proof>

thm *less-mod-eq-imp-add-divisor-le*

<proof>

lemma *iMOD-iprev'*: $t \in [r, \text{mod } m] \implies \text{iprev } (t + m) [r, \text{mod } m] = t$

<proof>

lemma *iMODb-inext*:

$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies$

$\text{inext } t [r, \text{mod } m, c] = t + m$

<proof>

lemma *iMODb-iprev*:

$\llbracket t \in [r, \text{mod } m, c]; r < t \rrbracket \implies$

$\text{iprev } t [r, \text{mod } m, c] = t - m$

<proof>

thm *mod-eq-imp-diff-mod-eq-divisor*

<proof>

thm *less-mod-eq-imp-add-divisor-le*

<proof>

lemma *iMODb-iprev'*:

$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies$

$\text{iprev } (t + m) [r, \text{mod } m, c] = t$

<proof>

thm *less-mod-eq-imp-add-divisor-le*

<proof>

lemmas *iT-inext* =

iFROM-inext

iTILL-inext

iIN-inext

iMOD-inext

iMODb-inext

lemmas *iT-iprev* =

iFROM-iprev'

iFROM-iprev

iTILL-iprev

iIN-iprev

iIN-iprev'

iMOD-iprev

iMOD-iprev'

iMODb-iprev

iMODb-iprev'

thm *iT-inext*

thm *iT-iprev*

thm *iprev-iMin*

thm *iT-finite[THEN inext-Max]*

lemma *iFROM-inext-if*:

$$\text{inext } t [n\dots] = (\text{if } t \in [n\dots] \text{ then Suc } t \text{ else } t)$$

<proof>

lemma *iTILL-inext-if*:

$$\text{inext } t [..n] = (\text{if } t < n \text{ then Suc } t \text{ else } t)$$

<proof>

lemma *iIN-inext-if*:

$$\text{inext } t [n\dots,d] = (\text{if } n \leq t \wedge t < n + d \text{ then Suc } t \text{ else } t)$$

<proof>

lemma *iMOD-inext-if*:

$$\text{inext } t [r, \text{mod } m] = (\text{if } t \in [r, \text{mod } m] \text{ then } t + m \text{ else } t)$$

<proof>

lemma *iMODb-inext-if*:

$$\text{inext } t [r, \text{mod } m, c] =$$

$$(\text{if } t \in [r, \text{mod } m, c] \wedge t < r + m * c \text{ then } t + m \text{ else } t)$$

<proof>

lemmas *iT-inext-if =*

iFROM-inext-if

iTILL-inext-if

iIN-inext-if

iMOD-inext-if

iMODb-inext-if

thm *iT-inext-if*

lemma *iFROM-iprev-if*:

$$\text{iprev } t [n\dots] = (\text{if } n < t \text{ then } t - \text{Suc } 0 \text{ else } t)$$

<proof>

lemma *iTILL-iprev-if*:

$$\text{iprev } t [..n] = (\text{if } t \in [..n] \text{ then } t - \text{Suc } 0 \text{ else } t)$$

<proof>

lemma *iIN-iprev-if*:

$$\text{iprev } t [n\dots,d] = (\text{if } n < t \wedge t \leq n + d \text{ then } t - \text{Suc } 0 \text{ else } t)$$

<proof>

lemma *iMOD-iprev-if*:

$$\text{iprev } t [r, \text{mod } m] =$$

$$(\text{if } t \in [r, \text{mod } m] \wedge r < t \text{ then } t - m \text{ else } t)$$

<proof>

lemma *iMODb-iprev-if*:

$$\text{iprev } t [r, \text{mod } m, c] =$$

$$(\text{if } t \in [r, \text{mod } m, c] \wedge r < t \text{ then } t - m \text{ else } t)$$

<proof>

lemmas *iT-iprev-if =*

iFROM-iprev-if

iTILL-iprev-if

iIN-iprev-if

iMOD-iprev-if
iMODb-iprev-if

thm *iT-iprev-if*

The difference between an element and the next/previous element is constant if the element is different from Min/Max of the interval

lemma *iFROM-inext-diff-const:*

$$t \in [n..] \implies \text{inext } t [n..] - t = \text{Suc } 0$$

<proof>

lemma *iFROM-iprev-diff-const:*

$$n < t \implies t - \text{iprev } t [n..] = \text{Suc } 0$$

<proof>

lemma *iFROM-iprev-diff-const':*

$$t \in [n..] \implies \text{Suc } t - \text{iprev } (\text{Suc } t) [n..] = \text{Suc } 0$$

<proof>

lemma *iTILL-inext-diff-const:*

$$t < n \implies \text{inext } t [..n] - t = \text{Suc } 0$$

<proof>

lemma *iTILL-iprev-diff-const:*

$$\llbracket t \in [..n]; 0 < t \rrbracket \implies t - \text{iprev } t [..n] = \text{Suc } 0$$

<proof>

thm *iIN-inext*

lemma *iIN-inext-diff-const:*

$$\llbracket n \leq t; t < n + d \rrbracket \implies \text{inext } t [n..,d] - t = \text{Suc } 0$$

<proof>

thm *iIN-iprev*

lemma *iIN-iprev-diff-const:*

$$\llbracket n < t; t \leq n + d \rrbracket \implies t - \text{iprev } t [n..,d] = \text{Suc } 0$$

<proof>

lemma *iIN-iprev-diff-const':*

$$\llbracket n \leq t; t < n + d \rrbracket \implies \text{Suc } t - \text{iprev } (\text{Suc } t) [n..,d] = \text{Suc } 0$$

<proof>

thm *iMOD-inext*

lemma *iMOD-inext-diff-const:*

$$t \in [r, \text{mod } m] \implies \text{inext } t [r, \text{mod } m] - t = m$$

<proof>

lemma *iMOD-iprev-diff-const':*

$$t \in [r, \text{mod } m] \implies (t + m) - \text{iprev } (t + m) [r, \text{mod } m] = m$$

<proof>

thm *iMOD-iprev*

lemma *iMOD-iprev-diff-const:*

$$\llbracket t \in [r, \text{mod } m]; r < t \rrbracket \implies t - \text{iprev } t [r, \text{mod } m] = m$$

<proof>

thm *iMODb-inext*

lemma *iMODb-inext-diff-const*:

$$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies \text{inext } t [r, \text{mod } m, c] - t = m$$

<proof>

thm *iMODb-iprev'*

lemma *iMODb-iprev-diff-const'*:

$$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies (t + m) - \text{iprev } (t + m) [r, \text{mod } m, c] = m$$

<proof>

thm *iMODb-iprev*

lemma *iMODb-iprev-diff-const*:

$$\llbracket t \in [r, \text{mod } m, c]; r < t \rrbracket \implies t - \text{iprev } t [r, \text{mod } m, c] = m$$

<proof>

lemmas *iT-inext-diff-const =*

iFROM-inext-diff-const

iTILL-inext-diff-const

iIN-inext-diff-const

iMOD-inext-diff-const

iMODb-inext-diff-const

lemmas *iT-iprev-diff-const =*

iFROM-iprev-diff-const

iFROM-iprev-diff-const'

iTILL-iprev-diff-const

iIN-iprev-diff-const

iIN-iprev-diff-const'

iMOD-iprev-diff-const'

iMOD-iprev-diff-const

iMODb-iprev-diff-const'

iMODb-iprev-diff-const

thm *iT-inext-diff-const*

thm *iT-iprev-diff-const*

1.7.1 Mirroring of intervals

thm *mirror-elem-def*

lemma

iIN-mirror-elem: *mirror-elem* $x [n \dots, d] = n + n + d - x$ **and**

iTILL-mirror-elem: *mirror-elem* $x [\dots, n] = n - x$ **and**

iMODb-mirror-elem: *mirror-elem* $x [r, \text{mod } m, c] = r + r + m * c - x$

<proof>

lemma *iMODb-imirror-bounds*:

$$r' + m' * c' \leq l + r \implies$$

$$\text{imirror-bounds } [r', \text{mod } m', c'] \text{ } l \text{ } r = [l + r - r' - m' * c', \text{mod } m', c']$$

<proof>

thm *mod-diff-right-eq*

<proof>

thm *mod-diff-right-eq*
 ⟨proof⟩
thm *mod-sub-eq-mod-swap*
 ⟨proof⟩
thm *mod-diff-right-eq*
 ⟨proof⟩

thm *imirror-bounds-def*
lemma *iIN-imirror-bounds*:
 $n + d \leq l + r \implies \text{imirror-bounds } [n\dots,d] \ l \ r = [l + r - n - d\dots,d]$
 ⟨proof⟩

lemma *iTILL-imirror-bounds*:
 $n \leq l + r \implies \text{imirror-bounds } [\dots,n] \ l \ r = [l + r - n\dots,n]$
 ⟨proof⟩

lemmas *iT-imirror-bounds =*
iTILL-imirror-bounds
iIN-imirror-bounds
iMODb-imirror-bounds
thm *iT-imirror-bounds*

lemma *iMODb-imirror-ident*: $\text{imirror } [r, \text{mod } m, c] = [r, \text{mod } m, c]$
 ⟨proof⟩
lemma *iIN-imirror-ident*: $\text{imirror } [n\dots,d] = [n\dots,d]$
 ⟨proof⟩
lemma *iTILL-imirror-ident*: $\text{imirror } [\dots,n] = [\dots,n]$
 ⟨proof⟩

lemmas *iT-imirror-ident =*
iTILL-imirror-ident
iIN-imirror-ident
iMODb-imirror-ident
thm *iT-imirror-ident*

1.7.2 Functions *inext-nth* and *iprev-nth* on intervals

lemma *iFROM-inext-nth* : $[n\dots] \rightarrow a = n + a$
 ⟨proof⟩
lemma *iIN-inext-nth* : $a \leq d \implies [n\dots,d] \rightarrow a = n + a$
 ⟨proof⟩
lemma *iIN-iprev-nth*: $a \leq d \implies [n\dots,d] \leftarrow a = n + d - a$
 ⟨proof⟩
lemma *iIN-inext-nth-if* :
 $[n\dots,d] \rightarrow a = (\text{if } a \leq d \text{ then } n + a \text{ else } n + d)$
 ⟨proof⟩
lemma *iIN-iprev-nth-if*:
 $[n\dots,d] \leftarrow a = (\text{if } a \leq d \text{ then } n + d - a \text{ else } n)$

<proof>

lemma *iTILL-inext-nth* : $a \leq n \implies [\dots n] \rightarrow a = a$

<proof>

lemma *iTILL-inext-nth-if* :

$[\dots n] \rightarrow a = (\text{if } a \leq n \text{ then } a \text{ else } n)$

<proof>

lemma *iTILL-iprev-nth*: $a \leq n \implies [\dots n] \leftarrow a = n - a$

<proof>

lemma *iTILL-iprev-nth-if*:

$[\dots n] \leftarrow a = (\text{if } a \leq n \text{ then } n - a \text{ else } 0)$

<proof>

lemma *iMOD-inext-nth*: $[r, \text{mod } m] \rightarrow a = r + m * a$

<proof>

lemma *iMODb-inext-nth*: $a \leq c \implies [r, \text{mod } m, c] \rightarrow a = r + m * a$

<proof>

lemma *iMODb-inext-nth-if*:

$[r, \text{mod } m, c] \rightarrow a = (\text{if } a \leq c \text{ then } r + m * a \text{ else } r + m * c)$

<proof>

lemma *iMODb-iprev-nth*:

$a \leq c \implies [r, \text{mod } m, c] \leftarrow a = r + m * (c - a)$

<proof>

lemma *iMODb-iprev-nth-if*:

$[r, \text{mod } m, c] \leftarrow a = (\text{if } a \leq c \text{ then } r + m * (c - a) \text{ else } r)$

<proof>

lemma *iIN-iFROM-inext-nth*:

$a \leq d \implies [n \dots, d] \rightarrow a = [n \dots] \rightarrow a$

<proof>

lemma *iIN-iFROM-inext*:

$a < n + d \implies \text{inext } a [n \dots, d] = \text{inext } a [n \dots]$

<proof>

lemma *iMOD-iMODb-inext-nth*:

$a \leq c \implies [r, \text{mod } m, c] \rightarrow a = [r, \text{mod } m] \rightarrow a$

<proof>

lemma *iMOD-iMODb-inext*:

$a < r + m * c \implies \text{inext } a [r, \text{mod } m, c] = \text{inext } a [r, \text{mod } m]$

<proof>

lemma *iMOD-inext-nth-Suc-diff*:

$([r, \text{mod } m] \rightarrow (\text{Suc } n)) - ([r, \text{mod } m] \rightarrow n) = m$

<proof>

lemma *iMOD-inext-nth-diff*:

$$([r, \text{mod } m] \rightarrow a) - ([r, \text{mod } m] \rightarrow b) = (a - b) * m$$

<proof>

lemma *iMODb-inext-nth-diff*: $\llbracket a \leq c; b \leq c \rrbracket \implies$
 $([r, \text{mod } m, c] \rightarrow a) - ([r, \text{mod } m, c] \rightarrow b) = (a - b) * m$
<proof>

1.8 Induction with intervals

thm

inext-induct

iFROM-inext

lemma *iFROM-induct*:

$$\llbracket P n; \bigwedge k. \llbracket k \in [n..]; P k \rrbracket \implies P (\text{Suc } k); a \in [n..] \rrbracket \implies P a$$

<proof>

lemma *iN-induct*:

$$\llbracket P n; \bigwedge k. \llbracket k \in [n..,d]; k \neq n + d; P k \rrbracket \implies P (\text{Suc } k); a \in [n..,d] \rrbracket \implies P a$$

<proof>

lemma *iTILL-induct*:

$$\llbracket P 0; \bigwedge k. \llbracket k \in [..n]; k \neq n; P k \rrbracket \implies P (\text{Suc } k); a \in [..n] \rrbracket \implies P a$$

<proof>

lemma *iMOD-induct*:

$$\llbracket P r; \bigwedge k. \llbracket k \in [r, \text{mod } m]; P k \rrbracket \implies P (k + m); a \in [r, \text{mod } m] \rrbracket \implies P a$$

<proof>

lemma *iMODb-induct*:

$$\llbracket P r; \bigwedge k. \llbracket k \in [r, \text{mod } m, c]; k \neq r + m * c; P k \rrbracket \implies P (k + m); a \in [r, \text{mod } m, c] \rrbracket \implies P a$$

<proof>

thm

iprev-induct

iN-inext

lemma *iN-rev-induct*:

$$\llbracket P (n + d); \bigwedge k. \llbracket k \in [n..,d]; k \neq n; P k \rrbracket \implies P (k - \text{Suc } 0); a \in [n..,d] \rrbracket \implies P a$$

<proof>

lemma *iTILL-rev-induct*:

$$\llbracket P n; \bigwedge k. \llbracket k \in [..n]; 0 < k; P k \rrbracket \implies P (k - \text{Suc } 0); a \in [..n] \rrbracket \implies P a$$

<proof>

lemma *iMODb-rev-induct*:

$$\llbracket P (r + m * c); \bigwedge k. \llbracket k \in [r, \text{mod } m, c]; k \neq r; P k \rrbracket \implies P (k - m); a \in [r, \text{mod } m, c] \rrbracket \implies P a$$

<proof>

end

2 IL-IntervalOperators: Arithmetic operators on natural intervals

theory *IL-IntervalOperators*
 imports *IL-Interval*
 begin

2.1 Arithmetic operations with intervals

2.1.1 Addition of and multiplication by constants

definition *iT-Plus* :: *iT* \Rightarrow *Time* \Rightarrow *iT* (**infixl** \oplus 55) where
 $I \oplus k \equiv (\lambda n. (n + k)) \text{ ' } I$

definition *iT-Mult* :: *iT* \Rightarrow *Time* \Rightarrow *iT* (**infixl** \otimes 55) where
 $iT\text{-Mult-def} : I \otimes k \equiv (\lambda n. (n * k)) \text{ ' } I$

lemma *iT-Plus-image-conv*: $I \oplus k = (\lambda n. (n + k)) \text{ ' } I$
 $\langle proof \rangle$

lemma *iT-Mult-image-conv*: $I \otimes k = (\lambda n. (n * k)) \text{ ' } I$
 $\langle proof \rangle$

lemma *iT-Plus-empty*: $\{\} \oplus k = \{\}$
 $\langle proof \rangle$

lemma *iT-Mult-empty*: $\{\} \otimes k = \{\}$
 $\langle proof \rangle$

lemma *iT-Plus-not-empty*: $I \neq \{\} \implies I \oplus k \neq \{\}$
 $\langle proof \rangle$

lemma *iT-Mult-not-empty*: $I \neq \{\} \implies I \otimes k \neq \{\}$
 $\langle proof \rangle$

lemma *iT-Plus-empty-iff*: $(I \oplus k = \{\}) = (I = \{\})$
 $\langle proof \rangle$

lemma *iT-Mult-empty-iff*: $(I \otimes k = \{\}) = (I = \{\})$
 $\langle proof \rangle$

lemma *iT-Plus-mono*: $A \subseteq B \implies A \oplus k \subseteq B \oplus k$
 $\langle proof \rangle$

lemma *iT-Mult-mono*: $A \subseteq B \implies A \otimes k \subseteq B \otimes k$
 $\langle proof \rangle$

lemma *iT-Mult-0*: $I \neq \{\} \implies I \otimes 0 = [\dots 0]$

<proof>

corollary *iT-Mult-0-if*: $I \otimes 0 = (\text{if } I = \{\} \text{ then } \{\} \text{ else } [\dots 0])$

<proof>

lemma *iT-Plus-mem-iff*: $x \in (I \oplus k) = (k \leq x \wedge (x - k) \in I)$

<proof>

lemma *iT-Plus-mem-iff2*: $x + k \in (I \oplus k) = (x \in I)$

<proof>

lemma *iT-Mult-mem-iff-0*: $x \in (I \otimes 0) = (I \neq \{\} \wedge x = 0)$

<proof>

lemma *iT-Mult-mem-iff*:

$0 < k \implies x \in (I \otimes k) = (x \text{ mod } k = 0 \wedge x \text{ div } k \in I)$

<proof>

lemma *iT-Mult-mem-iff2*: $0 < k \implies x * k \in (I \otimes k) = (x \in I)$

<proof>

lemma *iT-Plus-singleton*: $\{a\} \oplus k = \{a + k\}$

<proof>

lemma *iT-Mult-singleton*: $\{a\} \otimes k = \{a * k\}$

<proof>

lemma *iT-Plus-Un*: $(A \cup B) \oplus k = (A \oplus k) \cup (B \oplus k)$

<proof>

lemma *iT-Mult-Un*: $(A \cup B) \otimes k = (A \otimes k) \cup (B \otimes k)$

<proof>

lemma *iT-Plus-Int*: $(A \cap B) \oplus k = (A \oplus k) \cap (B \oplus k)$

<proof>

lemma *iT-Mult-Int*: $0 < k \implies (A \cap B) \otimes k = (A \otimes k) \cap (B \otimes k)$

<proof>

thm

iT-Plus-Un

iT-Mult-Un

iT-Plus-Int

iT-Mult-Int

thm *imageI*

lemma *iT-Plus-image*: $f \text{ ‘ } I \oplus n = (\lambda x. f x + n) \text{ ‘ } I$

<proof>

lemma *iT-Mult-image*: $f \text{ ‘ } I \otimes n = (\lambda x. f x * n) \text{ ‘ } I$

<proof>

lemma *iT-Plus-commute*: $I \oplus a \oplus b = I \oplus b \oplus a$
 ⟨proof⟩
lemma *iT-Mult-commute*: $I \otimes a \otimes b = I \otimes b \otimes a$
 ⟨proof⟩
lemma *iT-Plus-assoc*: $I \oplus a \oplus b = I \oplus (a + b)$
 ⟨proof⟩
lemma *iT-Mult-assoc*: $I \otimes a \otimes b = I \otimes (a * b)$
 ⟨proof⟩
lemma *iT-Plus-Mult-distrib*: $I \oplus n \otimes m = I \otimes m \oplus n * m$
 ⟨proof⟩⟨proof⟩⟨proof⟩⟨proof⟩⟨proof⟩⟨proof⟩
lemma *iT-Plus-finite-iff*: $\text{finite } (I \oplus k) = \text{finite } I$
 ⟨proof⟩
lemma *iT-Mult-0-finite*: $\text{finite } (I \otimes 0)$
 ⟨proof⟩
lemma *iT-Mult-finite-iff*: $0 < k \implies \text{finite } (I \otimes k) = \text{finite } I$
 ⟨proof⟩

lemma *iT-Plus-Min*: $I \neq \{\} \implies iMin (I \oplus k) = iMin I + k$
 ⟨proof⟩
lemma *iT-Mult-Min*: $I \neq \{\} \implies iMin (I \otimes k) = iMin I * k$
 ⟨proof⟩

thm *Max-mono2*

lemma *iT-Plus-Max*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies Max (I \oplus k) = Max I + k$
 ⟨proof⟩
lemma *iT-Mult-Max*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies Max (I \otimes k) = Max I * k$
 ⟨proof⟩

thm *iT-Mult-0*

corollary

iMOD-mult-0: $[r, \text{mod } m] \otimes 0 = [\dots 0]$ **and**
iMODb-mult-0: $[r, \text{mod } m, c] \otimes 0 = [\dots 0]$ **and**
iFROM-mult-0: $[n\dots] \otimes 0 = [\dots 0]$ **and**
iIN-mult-0: $[n\dots, d] \otimes 0 = [\dots 0]$ **and**
iTILL-mult-0: $[\dots n] \otimes 0 = [\dots 0]$
 ⟨proof⟩

lemmas *iT-mult-0* =

iTILL-mult-0
iFROM-mult-0
iIN-mult-0
iMOD-mult-0
iMODb-mult-0

thm *iT-mult-0*

lemma *iT-Plus-0*: $I \oplus 0 = I$
 ⟨proof⟩

lemma *iT-Mult-1*: $I \otimes \text{Suc } 0 = I$
 ⟨proof⟩

thm *iT-Plus-Min*

corollary

iFROM-add-Min: $iMin ([n..] \oplus k) = n + k$ **and**
iIN-add-Min: $iMin ([n..,d] \oplus k) = n + k$ **and**
iTILL-add-Min: $iMin ([..n] \oplus k) = k$ **and**
iMOD-add-Min: $iMin ([r, \text{mod } m] \oplus k) = r + k$ **and**
iMODb-add-Min: $iMin ([r, \text{mod } m, c] \oplus k) = r + k$
 ⟨proof⟩

corollary

iFROM-mult-Min: $iMin ([n..] \otimes k) = n * k$ **and**
iIN-mult-Min: $iMin ([n..,d] \otimes k) = n * k$ **and**
iTILL-mult-Min: $iMin ([..n] \otimes k) = 0$ **and**
iMOD-mult-Min: $iMin ([r, \text{mod } m] \otimes k) = r * k$ **and**
iMODb-mult-Min: $iMin ([r, \text{mod } m, c] \otimes k) = r * k$
 ⟨proof⟩

lemmas *iT-add-Min* =

iIN-add-Min
iTILL-add-Min
iFROM-add-Min
iMOD-add-Min
iMODb-add-Min

thm *iT-add-Min*

lemmas *iT-mult-Min* =

iIN-mult-Min
iTILL-mult-Min
iFROM-mult-Min
iMOD-mult-Min
iMODb-mult-Min

thm *iT-mult-Min*

lemma *iFROM-add*: $[n..] \oplus k = [n+k..]$
 ⟨proof⟩

lemma *iIN-add*: $[n..,d] \oplus k = [n+k..,d]$
 ⟨proof⟩

lemma *iTILL-add*: $[..i] \oplus k = [k..,i]$
 ⟨proof⟩

lemma *iMOD-add*: $[r, \text{mod } m] \oplus k = [r + k, \text{mod } m]$

<proof>

thm *mod-add*

<proof>

thm *mod-sub-add*[of k x]

<proof>

lemma *iMODb-add*: $[r, \text{mod } m, c] \oplus k = [r + k, \text{mod } m, c]$

<proof>

lemmas *iT-add* =

iMOD-add

iMODb-add

iFROM-add

iIN-add

iTILL-add

iT-Plus-singleton

thm *iT-add*

thm

mod-mult-distrib

mod-mult-distrib2

thm

mod-eq-mult-distrib

mod-factor-imp-mod-0

mod-factor-div

mod-factor-div-mod

lemma *iFROM-mult*: $[n \dots] \otimes k = [n * k, \text{mod } k]$

<proof>

lemma *iIN-mult*: $[n \dots, d] \otimes k = [n * k, \text{mod } k, d]$

<proof>

lemma *iTILL-mult*: $[\dots n] \otimes k = [0, \text{mod } k, n]$

<proof>

lemma *iMOD-mult*: $[r, \text{mod } m] \otimes k = [r * k, \text{mod } m * k]$

<proof>

thm *mod-mult2-eq*

<proof>

lemma *iMODb-mult*:

$[r, \text{mod } m, c] \otimes k = [r * k, \text{mod } m * k, c]$

<proof>

thm *iMODb-iMOD-iTILL-conv*

<proof>

thm *Int-absorb2*

<proof>
thm *iMOD-iTILL-iMODb-conv*
<proof>

lemmas *iT-mult* =
iTILL-mult
iFROM-mult
iIN-mult
iMOD-mult
iMODb-mult
iT-Mult-singleton
thm
iT-mult
iT-mult-0

2.1.2 Some conversions between intervals using constant addition and multiplication

lemma *iFROM-conv*: $[n\dots] = UNIV \oplus n$
<proof>
lemma *iIN-conv*: $[n\dots,d] = [\dots d] \oplus n$
<proof>
lemma *iMOD-conv*: $[r, \text{mod } m] = [0\dots] \otimes m \oplus r$
<proof>
lemma *iMODb-conv*: $[r, \text{mod } m, c] = [\dots c] \otimes m \oplus r$
<proof>

Some examples showing the utility of *iMODb_conv*

lemma $[12, \text{mod } 10, 4] = \{12, 22, 32, 42, 52\}$
<proof>
lemma $[12, \text{mod } 10, 4] = \{12, 22, 32, 42, 52\}$
<proof>
lemma $[12, \text{mod } 10, 4] = \{12, 22, 32, 42, 52\}$
<proof>
lemma $[r, \text{mod } m, 4] = \{r, r+m, r+2*m, r+3*m, r+4*m\}$
<proof>
thm *image-Collect*
<proof>

lemma $[2, \text{mod } 10, 4] = \{2, 12, 22, 32, 42\}$
<proof>

2.1.3 Subtraction of constants

definition *iT-Plus-neg* :: $iT \Rightarrow Time \Rightarrow iT$ (**infixl** $\oplus -$ 55) where
 $I \oplus - k \equiv \{x. x + k \in I\}$

lemma *iT-Plus-neg-mem-iff*: $(x \in I \oplus - k) = (x + k \in I)$

<proof>

thm *iT-Plus-mem-iff2*

lemma *iT-Plus-neg-mem-iff2*: $k \leq x \implies (x - k \in I \oplus - k) = (x \in I)$
<proof>

lemma *iT-Plus-neg-image-conv*: $I \oplus - k = (\lambda n.(n - k)) \cdot (I \downarrow \geq k)$

<proof>

thm *image-eqI*

<proof>

lemma *iT-Plus-neg-cut-eq*: $t \leq k \implies (I \downarrow \geq t) \oplus - k = I \oplus - k$

<proof>

lemma *iT-Plus-neg-mono*: $A \subseteq B \implies A \oplus - k \subseteq B \oplus - k$

<proof>

lemma *iT-Plus-neg-empty*: $\{\} \oplus - k = \{\}$

<proof>

lemma *iT-Plus-neg-Max-less-empty*:

$\llbracket \text{finite } I; \text{Max } I < k \rrbracket \implies I \oplus - k = \{\}$
<proof>

lemma *iT-Plus-neg-not-empty-iff*: $(I \oplus - k \neq \{\}) = (\exists x \in I. k \leq x)$

<proof>

lemma *iT-Plus-neg-empty-iff*:

$(I \oplus - k = \{\}) = (I = \{\}) \vee (\text{finite } I \wedge \text{Max } I < k)$
<proof>

lemma *iT-Plus-neg-assoc*: $(I \oplus - a) \oplus - b = I \oplus - (a + b)$

<proof>

lemma *iT-Plus-neg-commute*: $I \oplus - a \oplus - b = I \oplus - b \oplus - a$

<proof>

lemma *iT-Plus-neg-0*: $I \oplus - 0 = I$

<proof>

thm *diff-add-assoc*

lemma *iT-Plus-Plus-neg-assoc*: $b \leq a \implies I \oplus a \oplus - b = I \oplus (a - b)$

<proof>

lemma *iT-Plus-Plus-neg-assoc2*: $a \leq b \implies I \oplus a \oplus - b = I \oplus - (b - a)$

<proof>

lemma *iT-Plus-neg-Plus-le-cut-eq*:

$a \leq b \implies (I \oplus - a) \oplus b = (I \downarrow \geq a) \oplus (b - a)$
<proof>

corollary *iT-Plus-neg-Plus-le-Min-eq*:

$\llbracket a \leq b; a \leq iMin I \rrbracket \implies (I \oplus - a) \oplus b = I \oplus (b - a)$
 ⟨proof⟩

lemma *iT-Plus-neg-Plus-ge-cut-eq*:

$b \leq a \implies (I \oplus - a) \oplus b = (I \downarrow \geq a) \oplus - (a - b)$
 ⟨proof⟩

corollary *iT-Plus-neg-Plus-ge-Min-eq*:

$\llbracket b \leq a; a \leq iMin I \rrbracket \implies (I \oplus - a) \oplus b = I \oplus - (a - b)$
 ⟨proof⟩

lemma *iT-Plus-neg-Mult-distrib*:

$0 < m \implies I \oplus - n \otimes m = I \otimes m \oplus - n * m$
 ⟨proof⟩

thm *le-add-diff-inverse2*

lemma *iT-Plus-neg-Plus-le-inverse*: $k \leq iMin I \implies I \oplus - k \oplus k = I$
 ⟨proof⟩

lemma *iT-Plus-neg-Plus-inverse*: $I \oplus - k \oplus k = I \downarrow \geq k$
 ⟨proof⟩

thm *diff-add-inverse2*

lemma *iT-Plus-Plus-neg-inverse*: $I \oplus k \oplus - k = I$
 ⟨proof⟩

lemma *iT-Plus-neg-Un*: $(A \cup B) \oplus - k = (A \oplus - k) \cup (B \oplus - k)$
 ⟨proof⟩

lemma *iT-Plus-neg-Int*: $(A \cap B) \oplus - k = (A \oplus - k) \cap (B \oplus - k)$
 ⟨proof⟩

lemma *iT-Plus-neg-Max-singleton*: $\llbracket finite I; I \neq \{\} \rrbracket \implies I \oplus - Max I = \{0\}$
 ⟨proof⟩

lemma *iT-Plus-neg-singleton*: $\{a\} \oplus - k = (if\ k \leq a\ then\ \{a - k\}\ else\ \{\})$
 ⟨proof⟩

corollary *iT-Plus-neg-singleton1*: $k \leq a \implies \{a\} \oplus - k = \{a - k\}$
 ⟨proof⟩

corollary *iT-Plus-neg-singleton2*: $a < k \implies \{a\} \oplus - k = \{\}$
 ⟨proof⟩

lemma *iT-Plus-neg-finite-iff*: $finite (I \oplus - k) = finite I$
 ⟨proof⟩

thm *inj-on-finite-image-iff*
 ⟨proof⟩

thm *nat-cut-ge-finite-iff*
 ⟨proof⟩

lemma *iT-Plus-neg-Min*:

$$I \oplus - k \neq \{\} \implies iMin (I \oplus - k) = iMin (I \downarrow_{\geq} k) - k$$

<proof>

lemma *iT-Plus-neg-Max*:

$$\llbracket \text{finite } I; I \oplus - k \neq \{\} \rrbracket \implies Max (I \oplus - k) = Max I - k$$

<proof>

Subtractions of constants from intervals

lemma *iFROM-add-neg*: $[n\dots] \oplus - k = [n - k\dots]$

<proof>

lemma *iTILL-add-neg*: $[\dots n] \oplus - k = (\text{if } k \leq n \text{ then } [\dots n - k] \text{ else } \{\})$

<proof>

lemma *iTILL-add-neg1*: $k \leq n \implies [\dots n] \oplus - k = [\dots n - k]$

<proof>

lemma *iTILL-add-neg2*: $n < k \implies [\dots n] \oplus - k = \{\}$

<proof>

lemma *iIN-add-neg*:

$$[n\dots, d] \oplus - k = (\text{if } k \leq n \text{ then } [n - k\dots, d] \text{ else if } k \leq n + d \text{ then } [\dots n + d - k] \text{ else } \{\})$$

<proof>

lemma *iIN-add-neg1*: $k \leq n \implies [n\dots, d] \oplus - k = [n - k\dots, d]$

<proof>

lemma *iIN-add-neg2*: $\llbracket n \leq k; k \leq n + d \rrbracket \implies [n\dots, d] \oplus - k = [\dots n + d - k]$

<proof>

lemma *iIN-add-neg3*: $n + d < k \implies [n\dots, d] \oplus - k = \{\}$

<proof>

lemma *iMOD-0-add-neg*: $[r, \text{mod } 0] \oplus - k = \{r\} \oplus - k$

<proof>

lemma *iMOD-gr0-add-neg*:

$$0 < m \implies [r, \text{mod } m] \oplus - k = (\text{if } k \leq r \text{ then } [r - k, \text{mod } m] \text{ else } [(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m, \text{mod } m])$$

<proof>

thm *mod-sub-add[of k r m] le-diff-conv*

<proof>

thm *eq-commute[of r mod m]*

thm *mod-add-eq-mod-conv[of m x k r]*

<proof>

lemma *iMOD-add-neg*:

$$[r, \text{mod } m] \oplus - k = ($$

if $k \leq r$ then $[r - k, \text{mod } m]$
 else if $0 < m$ then $[(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m, \text{mod } m]$ else $\{\}$)

<proof>

corollary *iMOD-add-neg1:*

$$k \leq r \implies [r, \text{mod } m] \oplus - k = [r - k, \text{mod } m]$$

<proof>

lemma *iMOD-add-neg2:*

$$\llbracket 0 < m; r < k \rrbracket \implies [r, \text{mod } m] \oplus - k = [(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m, \text{mod } m]$$

<proof>

lemma *iMODb-mod-0-add-neg:* $[r, \text{mod } 0, c] \oplus - k = \{r\} \oplus - k$

<proof>

lemma *iMODb-add-neg:*

$$[r, \text{mod } m, c] \oplus - k = ($$

if $k \leq r$ then $[r - k, \text{mod } m, c]$
 else
 if $k \leq r + m * c$ then
 $[(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m, \text{mod } m, (r + m * c - k) \text{ div } m]$
 else $\{\}$)

<proof>

thm *mod-diff1-eq[of k]*

<proof>

thm *sub-mod-div-eq-div*

<proof>

lemma *iMODb-add-neg':*

$$[r, \text{mod } m, c] \oplus - k = ($$

if $k \leq r$ then $[r - k, \text{mod } m, c]$
 else if $k \leq r + m * c$ then
 if $k \text{ mod } m \leq r \text{ mod } m$
 then $[r \text{ mod } m - k \text{ mod } m, \text{mod } m, c + r \text{ div } m - k \text{ div } m]$
 else $[m + r \text{ mod } m - k \text{ mod } m, \text{mod } m, c + r \text{ div } m - \text{Suc } (k \text{ div } m)]$
 else $\{\}$)

<proof>

corollary *iMODb-add-neg1:*

$$k \leq r \implies [r, \text{mod } m, c] \oplus - k = [r - k, \text{mod } m, c]$$

<proof>

corollary *iMODb-add-neg2:*

$$\llbracket r < k; k \leq r + m * c \rrbracket \implies$$

$$[r, \text{mod } m, c] \oplus - k =$$

$$[(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m, \text{mod } m, (r + m * c - k) \text{ div } m]$$

<proof>

corollary *iMODb-add-neg2-mod-le*:

$$\begin{aligned} & \llbracket r < k; k \leq r + m * c; k \bmod m \leq r \bmod m \rrbracket \implies \\ & [r, \bmod m, c] \oplus - k = \\ & [r \bmod m - k \bmod m, \bmod m, c + r \operatorname{div} m - k \operatorname{div} m] \end{aligned}$$

<proof>

corollary *iMODb-add-neg2-mod-less*:

$$\begin{aligned} & \llbracket r < k; k \leq r + m * c; r \bmod m < k \bmod m \rrbracket \implies \\ & [r, \bmod m, c] \oplus - k = \\ & [m + r \bmod m - k \bmod m, \bmod m, c + r \operatorname{div} m - \operatorname{Suc} (k \operatorname{div} m)] \end{aligned}$$

<proof>

lemma *iMODb-add-neg3*: $r + m * c < k \implies [r, \bmod m, c] \oplus - k = \{\}$

<proof>

lemmas *iT-add-neg* =

iFROM-add-neg

iIN-add-neg

iTILL-add-neg

iMOD-add-neg

iMODb-add-neg

iT-Plus-neg-singleton

thm *iT-add-neg*

2.1.4 Subtraction of intervals from constants

definition *iT-Minus* :: *Time* \Rightarrow *iT* \Rightarrow *iT* (**infixl** \ominus 55) **where**

$$k \ominus I \equiv \{x. x \leq k \wedge (k - x) \in I\}$$

lemma *iT-Minus-mem-iff*: $(x \in k \ominus I) = (x \leq k \wedge k - x \in I)$

<proof>

lemma *iT-Minus-mono*: $A \subseteq B \implies k \ominus A \subseteq k \ominus B$

<proof>

lemma *iT-Minus-image-conv*: $k \ominus I = (\lambda x. k - x) \text{ ` } (I \downarrow \leq k)$

<proof>

lemma *iT-Minus-cut-eq*: $k \leq t \implies k \ominus (I \downarrow \leq t) = k \ominus I$

<proof>

lemma *iT-Minus-Minus-cut-eq*: $k \ominus (k \ominus (I \downarrow \leq k)) = I \downarrow \leq k$

<proof>

lemma $10 \ominus [\dots 3] = [7 \dots, 3]$

<proof>

lemma *iT-Minus-empty*: $k \ominus \{\} = \{\}$

<proof>

lemma *iT-Minus-0*: $k \ominus \{0\} = \{k\}$

<proof>

lemma *iT-Minus-bound*: $x \in k \ominus I \implies x \leq k$

<proof>

lemma *iT-Minus-finite*: *finite* ($k \ominus I$)

thm *finite-nat-iff-bounded-le2*

<proof>

lemma *iT-Minus-less-Min-empty*: $k < iMin\ I \implies k \ominus I = \{\}$

<proof>

lemma *iT-Minus-Min-singleton*: $I \neq \{\} \implies (iMin\ I) \ominus I = \{0\}$

<proof>

lemma *iT-Minus-empty-iff*: $(k \ominus I = \{\}) = (I = \{\} \vee k < iMin\ I)$

<proof>

lemma *iT-Minus-imirror-conv*:

$$k \ominus I = imirror\ (I \downarrow \leq k) \oplus k \oplus -\ (iMin\ I + Max\ (I \downarrow \leq k))$$

<proof>

thm *nat-add-right-cancel*[*THEN iffD2, of - - k*]

<proof>

thm *add-diff-assoc2*

<proof>

thm *iT-Minus-imirror-conv*

lemma *iT-Minus-imirror-conv'*:

$$k \ominus I = imirror\ (I \downarrow \leq k) \oplus k \oplus -\ (iMin\ (I \downarrow \leq k) + Max\ (I \downarrow \leq k))$$

<proof>

thm *cut-le-Min-eq*

<proof>

thm *iT-Minus-bound*

lemma *iT-Minus-Max*:

$$\llbracket I \neq \{\}; iMin\ I \leq k \rrbracket \implies Max\ (k \ominus I) = k - (iMin\ I)$$

thm *Max-equality*

<proof>

lemma *iT-Minus-Min*:

$$\llbracket I \neq \{\}; iMin\ I \leq k \rrbracket \implies iMin\ (k \ominus I) = k - (Max\ (I \downarrow \leq k))$$

<proof>

thm *subsetD*[*OF cut-le-subset, OF Max-in*]

<proof>

lemma *iT-Minus-Minus-eq*: $\llbracket finite\ I; Max\ I \leq k \rrbracket \implies k \ominus (k \ominus I) = I$

thm *iT-Minus-cut-eq*[*of k k I, symmetric*] *iT-Minus-Minus-cut-eq*

<proof>

lemma *iT-Minus-Minus-eq2*: $I \subseteq [\dots k] \implies k \ominus (k \ominus I) = I$
 ⟨proof⟩

thm *Max-subset*
 ⟨proof⟩

lemma *iT-Minus-Minus*: $a \ominus (b \ominus I) = (I \downarrow \leq b) \oplus a \oplus - b$
 ⟨proof⟩

lemma *iT-Minus-Plus-empty*: $k < n \implies k \ominus (I \oplus n) = \{\}$
 ⟨proof⟩

lemma *iT-Minus-Plus-commute*: $n \leq k \implies k \ominus (I \oplus n) = (k - n) \ominus I$
 ⟨proof⟩

lemma *iT-Minus-Plus-cut-assoc*: $(k \ominus I) \oplus n = (k + n) \ominus (I \downarrow \leq k)$
 ⟨proof⟩

lemma *iT-Minus-Plus-assoc*:

[[*finite* I ; $\text{Max } I \leq k$]] $\implies (k \ominus I) \oplus n = (k + n) \ominus I$
 ⟨proof⟩

lemma *iT-Minus-Plus-assoc2*:

$I \subseteq [\dots k] \implies (k \ominus I) \oplus n = (k + n) \ominus I$
 ⟨proof⟩

thm *Max-subset*
 ⟨proof⟩

lemma *iT-Minus-Un*: $k \ominus (A \cup B) = (k \ominus A) \cup (k \ominus B)$
 ⟨proof⟩

lemma *iT-Minus-Int*: $k \ominus (A \cap B) = (k \ominus A) \cap (k \ominus B)$
 ⟨proof⟩

lemma *iT-Minus-singleton*: $k \ominus \{a\} = (\text{if } a \leq k \text{ then } \{k - a\} \text{ else } \{\})$
 ⟨proof⟩

corollary *iT-Minus-singleton1*: $a \leq k \implies k \ominus \{a\} = \{k - a\}$
 ⟨proof⟩

corollary *iT-Minus-singleton2*: $k < a \implies k \ominus \{a\} = \{\}$
 ⟨proof⟩

lemma *iMOD-sub*:

$k \ominus [r, \text{mod } m] =$
 ($\text{if } r \leq k \text{ then } [(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m] \text{ else } \{\}$)
 ⟨proof⟩

thm *mod-sub-eq-mod-swap*[of r k x m , *symmetric*]
 ⟨proof⟩

corollary *iMOD-sub1*:

$$r \leq k \implies k \ominus [r, \text{mod } m] = [(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m]$$

<proof>

corollary *iMOD-sub2*: $k < r \implies k \ominus [r, \text{mod } m] = \{\}$

thm *iT-Minus-less-Min-empty*

<proof>

lemma *iTILL-sub*: $k \ominus [\dots n] = (\text{if } n \leq k \text{ then } [k - n \dots, n] \text{ else } [\dots k])$

<proof>

corollary *iTILL-sub1*: $n \leq k \implies k \ominus [\dots n] = [k - n \dots, n]$

<proof>

corollary *iTILL-sub2*: $k \leq n \implies k \ominus [\dots n] = [\dots k]$

<proof>

lemma *iMODb-sub*:

$$k \ominus [r, \text{mod } m, c] = (\text{if } r + m * c \leq k \text{ then } [k - r - m * c, \text{mod } m, c] \text{ else } \text{if } r \leq k \text{ then } [(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m] \text{ else } \{\})$$

<proof>

thm *iMODb-iMOD-iTILL-conv*

<proof>

thm *iT-Minus-Int*

<proof>

thm *add-le-imp-le-diff2*

<proof>

thm *le-diff-conv2[OF mod-le-dividend]*

<proof>

thm *iMODb-iMOD-iIN-conv*

<proof>

thm *iIN-inter*

<proof>

corollary *iMODb-sub1*:

$$\llbracket r \leq k; k \leq r + m * c \rrbracket \implies k \ominus [r, \text{mod } m, c] = [(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m]$$

<proof>

corollary *iMODb-sub2*: $k < r \implies k \ominus [r, \text{mod } m, c] = \{\}$

thm *iT-Minus-less-Min-empty*

<proof>

corollary *iMODb-sub3*:

$$r + m * c \leq k \implies k \ominus [r, \text{mod } m, c] = [k - r - m * c, \text{mod } m, c]$$

<proof>

lemma *iFROM-sub*: $k \ominus [n \dots] = (\text{if } n \leq k \text{ then } [\dots k - n] \text{ else } \{\})$

<proof>

corollary *iFROM-sub1*: $n \leq k \implies k \ominus [n \dots] = [\dots k - n]$

<proof>

corollary *iFROM-sub-empty*: $k < n \implies k \ominus [n \dots] = \{\}$

<proof>

lemma *iIN-sub*:

$k \ominus [n..d] = ($
if $n + d \leq k$ *then* $[k - (n + d)..d]$
else if $n \leq k$ *then* $[..k - n]$ *else* $\{\}$)

<proof>

lemma *iIN-sub1*: $n + d \leq k \implies k \ominus [n..d] = [k - (n + d)..d]$

<proof>

lemma *iIN-sub2*: $\llbracket n \leq k; k \leq n + d \rrbracket \implies k \ominus [n..d] = [..k - n]$

<proof>

lemma *iIN-sub3*: $k < n \implies k \ominus [n..d] = \{\}$

<proof>

lemmas *iT-sub* =

iFROM-sub

iIN-sub

iTILL-sub

iMOD-sub

iMODb-sub

iT-Minus-singleton

thm *iT-sub*

2.1.5 Division of intervals by constants

Monotonicity and injectivity of arithmetic operators

lemma *iMOD-div-right-strict-mono-on*:

$\llbracket 0 < k; k \leq m \rrbracket \implies \text{strict-mono-on } (\lambda x. x \text{ div } k) [r, \text{mod } m]$

<proof>

corollary *iMOD-div-right-inj-on*:

$\llbracket 0 < k; k \leq m \rrbracket \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{mod } m]$

<proof>

lemma *iMOD-mult-div-right-inj-on*:

inj-on $(\lambda x. x \text{ div } (k::\text{nat})) [r, \text{mod } (k * m)]$

<proof>

lemma *iMOD-mult-div-right-inj-on2*:

$m \text{ mod } k = 0 \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{mod } m]$

<proof>

lemma *iMODb-div-right-strict-mono-on*:

$\llbracket 0 < k; k \leq m \rrbracket \implies \text{strict-mono-on } (\lambda x. x \text{ div } k) [r, \text{mod } m, c]$

thm *strict-mono-on-subset[OF iMOD-div-right-strict-mono-on iMODb-iMOD-subset-same]*

<proof>

corollary *iMODb-div-right-inj-on*:

$\llbracket 0 < k; k \leq m \rrbracket \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{mod } m, c]$

<proof>

lemma *iMODb-mult-div-right-inj-on*:

inj-on $(\lambda x. x \text{ div } (k::\text{nat})) [r, \text{mod } (k * m), c]$

thm *subset-inj-on*[*OF iMOD-div-right-inj-on iMODb-iMOD-subset-same*]

<proof>

corollary *iMODb-mult-div-right-inj-on2*:

$m \text{ mod } k = 0 \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{mod } m, c]$

<proof>

definition *iT-Div* :: *iT* \Rightarrow *Time* \Rightarrow *iT* (**infixl** \circ 55) **where**

$I \circ k \equiv (\lambda n. (n \text{ div } k)) ' I$

lemma *iT-Div-image-conv*: $I \circ k = (\lambda n. (n \text{ div } k)) ' I$

<proof>

lemma *iT-Div-mono*: $A \subseteq B \implies A \circ k \subseteq B \circ k$

<proof>

lemma *iT-Div-empty*: $\{\} \circ k = \{\}$

<proof>

lemma *iT-Div-not-empty*: $I \neq \{\} \implies I \circ k \neq \{\}$

<proof>

lemma *iT-Div-empty-iff*: $(I \circ k = \{\}) = (I = \{\})$

<proof>

lemma *iT-Div-0*: $I \neq \{\} \implies I \circ 0 = [\dots 0]$

<proof>

corollary *iT-Div-0-if*: $I \circ 0 = (\text{if } I = \{\} \text{ then } \{\} \text{ else } [\dots 0])$

<proof>

corollary

iFROM-div-0: $[n\dots] \circ 0 = [\dots 0]$ **and**

iTILL-div-0: $[\dots n] \circ 0 = [\dots 0]$ **and**

iIN-div-0: $[n\dots, d] \circ 0 = [\dots 0]$ **and**

iMOD-div-0: $[r, \text{mod } m] \circ 0 = [\dots 0]$ **and**

iMODb-div-0: $[r, \text{mod } m, c] \circ 0 = [\dots 0]$

<proof>

lemmas *iT-div-0* =

iTILL-div-0

iFROM-div-0

iIN-div-0

iMOD-div-0

iMODb-div-0

thm *iT-div-0*

lemma *iT-Div-1*: $I \circ \text{Suc } 0 = I$

<proof>

lemma *iT-Div-mem-iff-0*: $x \in (I \circ 0) = (I \neq \{\}) \wedge x = 0$

<proof>

lemma *iT-Div-mem-iff*:

$0 < k \implies x \in (I \circ k) = (\exists y \in I. y \text{ div } k = x)$

<proof>

lemma *iT-Div-mem-iff2*:

$0 < k \implies x \text{ div } k \in (I \circ k) = (\exists y \in I. y \text{ div } k = x \text{ div } k)$

<proof>

lemma *iT-Div-mem-iff-Int*:

$0 < k \implies x \in (I \circ k) = (I \cap [x * k \dots k - \text{Suc } 0] \neq \{\})$

thm *ex-in-conv[symmetric]*

<proof>

thm *le-less-div-conv*

<proof>

thm *less-eq-le-pred*

<proof>

lemma *iT-Div-imp-mem*:

$0 < k \implies x \in I \implies x \text{ div } k \in (I \circ k)$

<proof>

lemma *iT-Div-singleton*: $\{a\} \circ k = \{a \text{ div } k\}$

<proof>

lemma *iT-Div-Un*: $(A \cup B) \circ k = (A \circ k) \cup (B \circ k)$

<proof>

lemma *iT-Div-insert*: $(\text{insert } n \ I) \circ k = \text{insert } (n \text{ div } k) \ (I \circ k)$

<proof>

lemma *not-iT-Div-Int*: $\neg (\forall k \ A \ B. (A \cap B) \circ k = (A \circ k) \cap (B \circ k))$

<proof>

thm *subset-image-Int*

lemma *subset-iT-Div-Int*: $A \subseteq B \implies (A \cap B) \circ k = (A \circ k) \cap (B \circ k)$

<proof>

lemma *iFROM-iT-Div-Int:*

$$\llbracket 0 < k; n \leq iMin A \rrbracket \implies (A \cap [n\dots]) \circledast k = (A \circledast k) \cap ([n\dots] \circledast k)$$

<proof>

lemma *iIN-iT-Div-Int:*

$$\llbracket 0 < k; n \leq iMin A; \forall x \in A. x \text{ div } k \leq (n + d) \text{ div } k \implies x \leq n + d \rrbracket \implies$$

$$(A \cap [n\dots, d]) \circledast k = (A \circledast k) \cap ([n\dots, d] \circledast k)$$

<proof>

corollary *iTILL-iT-Div-Int:*

$$\llbracket 0 < k; \forall x \in A. x \text{ div } k \leq n \text{ div } k \implies x \leq n \rrbracket \implies$$

$$(A \cap [\dots n]) \circledast k = (A \circledast k) \cap ([\dots n] \circledast k)$$

<proof>

lemma *iIN-iT-Div-Int-mod-0:*

$$\llbracket 0 < k; n \text{ mod } k = 0; \forall x \in A. x \text{ div } k \leq (n + d) \text{ div } k \implies x \leq n + d \rrbracket \implies$$

$$(A \cap [n\dots, d]) \circledast k = (A \circledast k) \cap ([n\dots, d] \circledast k)$$

<proof>

thm *div-le-mono*

<proof>

thm *less-mod-0-imp-div-less*

<proof>

lemma *mod-partition-iT-Div-Int:*

$$\llbracket 0 < k; 0 < d \rrbracket \implies$$

$$(A \cap [n * k\dots, d * k - Suc 0]) \circledast k =$$

$$(A \circledast k) \cap ([n * k\dots, d * k - Suc 0] \circledast k)$$

thm *iIN-iT-Div-Int-mod-0*

*<proof>**<proof>*

corollary *mod-partition-iT-Div-Int2:*

$$\llbracket 0 < k; 0 < d; n \text{ mod } k = 0; d \text{ mod } k = 0 \rrbracket \implies$$

$$(A \cap [n\dots, d - Suc 0]) \circledast k =$$

$$(A \circledast k) \cap ([n\dots, d - Suc 0] \circledast k)$$

<proof>

thm *mod-partition-iT-Div-Int*

<proof>

corollary *mod-partition-iT-Div-Int-one-segment:*

$$0 < k \implies$$

$$(A \cap [n * k\dots, k - Suc 0]) \circledast k = (A \circledast k) \cap ([n * k\dots, k - Suc 0] \circledast k)$$

<proof>

corollary *mod-partition-iT-Div-Int-one-segment2:*

$$\llbracket 0 < k; n \text{ mod } k = 0 \rrbracket \implies$$

$$(A \cap [n\dots, k - Suc 0]) \circledast k = (A \circledast k) \cap ([n\dots, k - Suc 0] \circledast k)$$

<proof>

thm

iT-Div-Un

subset-iT-Div-Int

thm

mod-partition-iT-Div-Int
mod-partition-iT-Div-Int2

lemma *iT-Div-assoc*: $I \odot a \odot b = I \odot (a * b)$
 ⟨proof⟩

lemma *iT-Div-commute*: $I \odot a \odot b = I \odot b \odot a$
 ⟨proof⟩

lemma *iT-Mult-Div-self*: $0 < k \implies I \otimes k \odot k = I$
 ⟨proof⟩

lemma *iT-Mult-Div*:
 $\llbracket 0 < d; k \bmod d = 0 \rrbracket \implies I \otimes k \odot d = I \otimes (k \text{ div } d)$
 ⟨proof⟩

lemma *iT-Div-Mult-self*:
 $0 < k \implies I \odot k \otimes k = \{y. \exists x \in I. y = x - x \bmod k\}$
 ⟨proof⟩

thm *div-add1-eq*

lemma *iT-Plus-Div-distrib-mod-less*:

$\forall x \in I. x \bmod m + n \bmod m < m \implies I \oplus n \odot m = I \odot m \oplus n \text{ div } m$
 ⟨proof⟩

corollary *iT-Plus-Div-distrib-mod-0*:

$n \bmod m = 0 \implies I \oplus n \odot m = I \odot m \oplus n \text{ div } m$
 ⟨proof⟩

lemma *iT-Div-Min*: $I \neq \{\}$ $\implies iMin (I \odot k) = iMin I \text{ div } k$
 ⟨proof⟩

thm *iT-Div-Min*

corollary

iFROM-div-Min: $iMin ([n..] \odot k) = n \text{ div } k$ **and**

iIN-div-Min: $iMin ([n..,d] \odot k) = n \text{ div } k$ **and**

iTILL-div-Min: $iMin ([..n] \odot k) = 0$ **and**

iMOD-div-Min: $iMin ([r, \bmod m] \odot k) = r \text{ div } k$ **and**

iMODb-div-Min: $iMin ([r, \bmod m, c] \odot k) = r \text{ div } k$

⟨proof⟩

lemmas *iT-div-Min* =

iFROM-div-Min

iIN-div-Min

iTILL-div-Min

iMOD-div-Min

iMODb-div-Min

thm *iT-div-Min*

lemma *iT-Div-Max*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies Max (I \odot k) = Max I \text{ div } k$
 ⟨proof⟩

corollary

iIN-div-Max: $Max ([n\dots,d] \odot k) = (n + d) \text{ div } k$ **and**

iTILL-div-Max: $Max ([\dots n] \odot k) = n \text{ div } k$ **and**

iMODb-div-Max: $Max ([r, \text{mod } m, c] \odot k) = (r + m * c) \text{ div } k$
 ⟨proof⟩

lemma *iT-Div-0-finite*: $finite (I \odot 0)$

⟨proof⟩

lemma *iT-Div-infinite-iff*: $0 < k \implies infinite (I \odot k) = infinite I$

⟨proof⟩

lemma *iT-Div-finite-iff*: $0 < k \implies finite (I \odot k) = finite I$

⟨proof⟩

lemma *iFROM-div*: $0 < k \implies [n\dots] \odot k = [n \text{ div } k\dots]$

⟨proof⟩

thm *div-add1-eq*

lemma *iIN-div*:

$0 < k \implies$

$[n\dots,d] \odot k = [n \text{ div } k\dots, d \text{ div } k + (n \text{ mod } k + d \text{ mod } k) \text{ div } k]$

⟨proof⟩

corollary *iIN-div-if*:

$0 < k \implies [n\dots,d] \odot k =$

$[n \text{ div } k\dots, d \text{ div } k + (\text{if } n \text{ mod } k + d \text{ mod } k < k \text{ then } 0 \text{ else } \text{Suc } 0)]$

⟨proof⟩

corollary *iIN-div-eq1*:

$\llbracket 0 < k; n \text{ mod } k + d \text{ mod } k < k \rrbracket \implies$

$[n\dots,d] \odot k = [n \text{ div } k\dots, d \text{ div } k]$

⟨proof⟩

corollary *iIN-div-eq2*:

$\llbracket 0 < k; k \leq n \text{ mod } k + d \text{ mod } k \rrbracket \implies$

$[n\dots,d] \odot k = [n \text{ div } k\dots, \text{Suc } (d \text{ div } k)]$

⟨proof⟩

corollary *iIN-div-mod-eq-0*:

$\llbracket 0 < k; n \text{ mod } k = 0 \rrbracket \implies [n\dots,d] \odot k = [n \text{ div } k\dots, d \text{ div } k]$

⟨proof⟩

lemma *iTILL-div*:

$0 < k \implies [\dots n] \odot k = [\dots n \text{ div } k]$

⟨proof⟩

lemma *iMOD-div-ge*:

$\llbracket 0 < m; m \leq k \rrbracket \implies [r, \text{mod } m] \odot k = [r \text{ div } k \dots]$
 ⟨proof⟩

thm *div-mult-cancel*

⟨proof⟩

corollary *iMOD-div-self*:

$0 < m \implies [r, \text{mod } m] \odot m = [r \text{ div } m \dots]$
 ⟨proof⟩

lemma *iMOD-div*:

$\llbracket 0 < k; m \text{ mod } k = 0 \rrbracket \implies$
 $[r, \text{mod } m] \odot k = [r \text{ div } k, \text{mod } (m \text{ div } k)]$
 ⟨proof⟩

thm *iMOD-mult[of r div k q k]*

⟨proof⟩

thm *arg-cong[where f=λx. x ⊕ (r mod k)]*

⟨proof⟩

thm *iT-Plus-Div-distrib-mod-less*

⟨proof⟩

thm *mod-factor-imp-mod-0[of x q k]*

⟨proof⟩

thm *iT-Mult-Div[OF - mod-self]*

⟨proof⟩

thm

iMOD-div

iMOD-div-ge

lemma *iMODb-div-self*:

$0 < m \implies [r, \text{mod } m, c] \odot m = [r \text{ div } m \dots, c]$

thm *iMODb-iMOD-iTILL-conv*

⟨proof⟩

thm *iTILL-iT-Div-Int*

⟨proof⟩

thm *div-le-mod-le-imp-le*

⟨proof⟩

lemma *iMODb-div-ge*:

$\llbracket 0 < m; m \leq k \rrbracket \implies$
 $[r, \text{mod } m, c] \odot k = [r \text{ div } k \dots, (r + m * c) \text{ div } k - r \text{ div } k]$
 ⟨proof⟩

thm *iMODb-append-union-Suc[of r m c 0, symmetric]*

⟨proof⟩

thm *div-add1-eq1*

⟨proof⟩

thm *div-add1-eq2*

⟨proof⟩

thm *iIN-Suc-insert-conv*

⟨proof⟩

thm *div-add1-eq-if*

corollary *iMODb-div-ge-if*:

$\llbracket 0 < m; m \leq k \rrbracket \implies$
 $[r, \text{mod } m, c] \odot k =$
 $[r \text{ div } k \dots, m * c \text{ div } k + (\text{if } r \text{ mod } k + m * c \text{ mod } k < k \text{ then } 0 \text{ else } \text{Suc } 0)]$
 ⟨proof⟩

thm *iMODb-div-ge*

lemma *iMODb-div*:

$\llbracket 0 < k; m \text{ mod } k = 0 \rrbracket \implies$
 $[r, \text{mod } m, c] \odot k = [r \text{ div } k, \text{mod } (m \text{ div } k), c]$
 ⟨proof⟩

thm *iTILL-iT-Div-Int*[of $k [r, \text{mod } m] m * c$]

⟨proof⟩

thm *mod-0-imp-mod-mult-right-0*

⟨proof⟩

thm *mod-eq-mod-0-imp-mod-eq*[of $x m r k$]

⟨proof⟩

lemmas *iT-div =*

iTILL-div

iFROM-div

iIN-div

iMOD-div

iMODb-div

iT-Div-singleton

thm

iT-div

iT-div-0

This lemma is valid for all $k \leq m$, i. e., also for k with $m \text{ mod } k \neq (0::'a)$.

lemma *iMODb-div-unique*:

$\llbracket 0 < k; k \leq m; k \leq c; [r', \text{mod } m', c'] = [r, \text{mod } m, c] \odot k \rrbracket \implies$
 $r' = r \text{ div } k \wedge m' = m \text{ div } k \wedge c' = c$
 ⟨proof⟩

thm *iT-Div-Min*

⟨proof⟩

thm *div-add1-eq-if*[of $k r m * c$]

⟨proof⟩

thm *card-image*[OF *iMODb-div-right-inj-on*]

⟨proof⟩

thm *iMODb-div-right-strict-mono-on*

⟨proof⟩

thm *iMODb-inext-nth-diff*

⟨proof⟩

thm *inext-nth-image*[OF *iMODb-not-empty*]

⟨proof⟩

lemma *iMODb-div-mod-gr0-is-0-not-ex0*:

$$\llbracket 0 < k; k < m; 0 < m \bmod k; k \leq c; r \bmod k = 0 \rrbracket \implies$$

$$\neg(\exists r' m' c'. [r', \bmod m', c'] = [r, \bmod m, c] \otimes k)$$

<proof>

thm *iMODb-div-unique*

<proof>

thm *div-mult1-eq[of m c]*

<proof>

lemma *iMODb-div-mod-gr0-not-ex-arith-aux1*:

$$\llbracket (0::nat) < k; k < m; 0 < x1 \rrbracket \implies$$

$$x1 * m + x2 - x \bmod k + x^3 + x \bmod k = x1 * m + x^2 + x^3$$

<proof>

lemma *iMODb-div-mod-gr0-not-ex*:

$$\llbracket 0 < k; k < m; 0 < m \bmod k; k \leq c \rrbracket \implies$$

$$\neg(\exists r' m' c'. [r', \bmod m', c'] = [r, \bmod m, c] \otimes k)$$

<proof>

thm *iMODb-div-mod-gr0-is-0-not-ex0*

<proof>

thm *iMODb-div-unique*

<proof>

thm *div-add1-eq*

<proof>

thm *div-mult1-eq[of c m k]*

<proof>

thm *iMODb-div-mod-gr0-not-ex*

thm

cut-le-image

iMOD-div-right-strict-mono-on

iMOD-cut-le

thm *card-image[OF strict-mono-on-imp-inj-on]*

lemma *iMOD-div-eq-imp-iMODb-div-eq*:

$$\llbracket 0 < k; k \leq m; [r', \bmod m'] = [r, \bmod m] \otimes k \rrbracket \implies$$

$$[r', \bmod m', c] = [r, \bmod m, c] \otimes k$$

<proof>

thm *iMODb-div-right-strict-mono-on[of k m r c]*

thm *card-image[OF strict-mono-on-imp-inj-on]*

thm *card-image[OF strict-mono-on-imp-inj-on[OF iMODb-div-right-strict-mono-on[of k m r c]]]*

<proof>

thm *iMOD-cut-le[of r m r + m * c]*

<proof>

thm *cut-le-image[OF - subset-refl]*

<proof>

thm *div-le-mono*

<proof>

lemma *iMOD-div-unique*:

$$\llbracket 0 < k; k \leq m; [r', \text{mod } m'] = [r, \text{mod } m] \circledast k \rrbracket \implies \\ r' = r \text{ div } k \wedge m' = m \text{ div } k$$

thm *iMOD-div-eq-imp-iMODb-div-eq*

<proof>

thm *iMODb-div-unique[of k - k]*

<proof>

thm *iMODb-div-mod-gr0-not-ex*

lemma *iMOD-div-mod-gr0-not-ex*:

$$\llbracket 0 < k; k < m; 0 < m \text{ mod } k \rrbracket \implies \\ \neg (\exists r' m'. [r', \text{mod } m'] = [r, \text{mod } m] \circledast k)$$

<proof>

thm *iMOD-div-eq-imp-iMODb-div-eq[OF - less-imp-le]*

<proof>

thm *iMODb-div-mod-gr0-not-ex[of k m k r]*

<proof>

2.2 Interval cut operators with arithmetic interval operators

lemma

$$iT\text{-Plus-cut-le2}: \quad (I \oplus k) \downarrow \leq (t + k) = (I \downarrow \leq t) \oplus k \text{ and}$$

$$iT\text{-Plus-cut-less2}: \quad (I \oplus k) \downarrow < (t + k) = (I \downarrow < t) \oplus k \text{ and}$$

$$iT\text{-Plus-cut-ge2}: \quad (I \oplus k) \downarrow \geq (t + k) = (I \downarrow \geq t) \oplus k \text{ and}$$

$$iT\text{-Plus-cut-greater2}: \quad (I \oplus k) \downarrow > (t + k) = (I \downarrow > t) \oplus k$$

<proof>

lemma *iT-Plus-cut-le*:

$$(I \oplus k) \downarrow \leq t = (\text{if } t < k \text{ then } \{\} \text{ else } I \downarrow \leq (t - k) \oplus k)$$

<proof>

thm *iT-Plus-cut-le2[of I k t - k]*

<proof>

lemma *iT-Plus-cut-less*: $(I \oplus k) \downarrow < t = I \downarrow < (t - k) \oplus k$

<proof>

lemma *iT-Plus-cut-ge*: $(I \oplus k) \downarrow \geq t = I \downarrow \geq (t - k) \oplus k$

<proof>

lemma *iT-Plus-cut-greater*:

$$(I \oplus k) \downarrow > t = (\text{if } t < k \text{ then } I \oplus k \text{ else } I \downarrow > (t - k) \oplus k)$$

<proof>

lemma

$$iT\text{-Mult-cut-le2}: \quad 0 < k \implies (I \otimes k) \downarrow \leq (t * k) = (I \downarrow \leq t) \otimes k \text{ and}$$

$$iT\text{-Mult-cut-less2}: \quad 0 < k \implies (I \otimes k) \downarrow < (t * k) = (I \downarrow < t) \otimes k \text{ and}$$

$$iT\text{-Mult-cut-ge2}: \quad 0 < k \implies (I \otimes k) \downarrow \geq (t * k) = (I \downarrow \geq t) \otimes k \text{ and}$$

$$iT\text{-Mult-cut-greater2}: \quad 0 < k \implies (I \otimes k) \downarrow > (t * k) = (I \downarrow > t) \otimes k$$

<proof>

lemma *iT-Mult-cut-le*:

$$0 < k \implies (I \otimes k) \downarrow \leq t = (I \downarrow \leq (t \text{ div } k)) \otimes k$$

<proof>

lemma *iT-Mult-cut-less*:

$$0 < k \implies (I \otimes k) \downarrow < t =$$

$$(if\ t\ mod\ k = 0\ then\ (I \downarrow < (t \text{ div } k))\ else\ I \downarrow < Suc\ (t \text{ div } k)) \otimes k$$

<proof>

thm *less-mod-0-imp-div-less*[of $t\ x\ k$]

<proof>

lemma *iT-Mult-cut-greater*:

$$0 < k \implies (I \otimes k) \downarrow > t = (I \downarrow > (t \text{ div } k)) \otimes k$$

<proof>

lemma *iT-Mult-cut-ge*:

$$0 < k \implies (I \otimes k) \downarrow \geq t =$$

$$(if\ t\ mod\ k = 0\ then\ (I \downarrow \geq (t \text{ div } k))\ else\ I \downarrow \geq Suc\ (t \text{ div } k)) \otimes k$$

<proof>

thm *div-le-mono*[OF *order-less-imp-le*]

<proof>

lemma *iT-Plus-neg-cut-le2*: $k \leq t \implies (I \oplus - k) \downarrow \leq (t - k) = (I \downarrow \leq t) \oplus - k$

<proof>

thm *i-cut-commute-disj*[of $op\ \downarrow \leq\ op\ \downarrow \geq$]

<proof>

lemma *iT-Plus-neg-cut-less2*: $(I \oplus - k) \downarrow < (t - k) = (I \downarrow < t) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-ge2*: $(I \oplus - k) \downarrow \geq (t - k) = (I \downarrow \geq t) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-greater2*: $k \leq t \implies (I \oplus - k) \downarrow > (t - k) = (I \downarrow > t) \oplus - k$

k

<proof>

lemma *iT-Plus-neg-cut-le*: $(I \oplus - k) \downarrow \leq t = I \downarrow \leq (t + k) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-less*: $(I \oplus - k) \downarrow < t = I \downarrow < (t + k) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-ge*: $(I \oplus - k) \downarrow \geq t = I \downarrow \geq (t + k) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-greater*: $(I \oplus - k) \downarrow > t = I \downarrow > (t + k) \oplus - k$

<proof>

lemma *iT-Minus-cut-le2*: $t \leq k \implies (k \ominus I) \downarrow \leq (k - t) = k \ominus (I \downarrow \geq t)$

<proof>

lemma *iT-Minus-cut-less2*: $(k \ominus I) \downarrow < (k - t) = k \ominus (I \downarrow > t)$

<proof>

lemma *iT-Minus-cut-ge2*: $(k \ominus I) \downarrow \geq (k - t) = k \ominus (I \downarrow \leq t)$

<proof>

lemma *iT-Minus-cut-greater2*: $t \leq k \implies (k \ominus I) \downarrow > (k - t) = k \ominus (I \downarrow < t)$

<proof>

lemma *iT-Minus-cut-le*: $(k \ominus I) \downarrow \leq t = k \ominus (I \downarrow \geq (k - t))$

<proof>

lemma *iT-Minus-cut-less*:

$(k \ominus I) \downarrow < t = (\text{if } t \leq k \text{ then } k \ominus (I \downarrow > (k - t)) \text{ else } k \ominus I)$

<proof>

lemma *iT-Minus-cut-ge*:

$(k \ominus I) \downarrow \geq t = (\text{if } t \leq k \text{ then } k \ominus (I \downarrow \leq (k - t)) \text{ else } \{\})$

<proof>

lemma *iT-Minus-cut-greater*: $(k \ominus I) \downarrow > t = k \ominus (I \downarrow < (k - t))$

<proof>

thm *iT-Div-def*

thm *iT-Mult-cut-le2*

thm *iT-Div-mem-iff*

lemma *iT-Div-cut-le*:

$0 < k \implies (I \otimes k) \downarrow \leq t = I \downarrow \leq (t * k + (k - \text{Suc } 0)) \otimes k$

<proof>

thm *div-le-conv*

<proof>

lemma *iT-Div-cut-less*:

$0 < k \implies (I \otimes k) \downarrow < t = I \downarrow < (t * k) \otimes k$

<proof>

lemma *iT-Div-cut-ge*:

$0 < k \implies (I \otimes k) \downarrow \geq t = I \downarrow \geq (t * k) \otimes k$

<proof>

thm *le-div-conv*

<proof>

lemma *iT-Div-cut-greater*:

$0 < k \implies (I \otimes k) \downarrow > t = I \downarrow > (t * k + (k - \text{Suc } 0)) \otimes k$

<proof>

lemma *iT-Div-cut-le2*:

$0 < k \implies (I \otimes k) \downarrow \leq (t \text{ div } k) = I \downarrow \leq (t - t \text{ mod } k + (k - \text{Suc } 0)) \otimes k$

<proof>

lemma *iT-Div-cut-less2*:

$0 < k \implies (I \otimes k) \downarrow < (t \text{ div } k) = I \downarrow < (t - t \text{ mod } k) \otimes k$

<proof>

lemma *iT-Div-cut-ge2*:

$0 < k \implies (I \otimes k) \downarrow \geq (t \text{ div } k) = I \downarrow \geq (t - t \text{ mod } k) \otimes k$

<proof>

lemma *iT-Div-cut-greater2*:

$0 < k \implies (I \otimes k) \downarrow > (t \text{ div } k) = I \downarrow > (t - t \text{ mod } k + (k - \text{Suc } 0)) \otimes k$
 ⟨proof⟩

2.3 *inext* and *iprev* with interval operators

lemma *iT-Plus-inext*: $\text{inext } (n + k) (I \oplus k) = (\text{inext } n I) + k$
 ⟨proof⟩

lemma *iT-Plus-iprev*: $\text{iprev } (n + k) (I \oplus k) = (\text{iprev } n I) + k$
 ⟨proof⟩

lemma *iT-Plus-inext2*: $k \leq n \implies \text{inext } n (I \oplus k) = (\text{inext } (n - k) I) + k$
 ⟨proof⟩

lemma *iT-Plus-prev2*: $k \leq n \implies \text{iprev } n (I \oplus k) = (\text{iprev } (n - k) I) + k$
 ⟨proof⟩

lemma *iT-Mult-inext*: $\text{inext } (n * k) (I \otimes k) = (\text{inext } n I) * k$
 ⟨proof⟩

lemma *iT-Mult-iprev*: $\text{iprev } (n * k) (I \otimes k) = (\text{iprev } n I) * k$
 ⟨proof⟩

lemma *iT-Mult-inext2-if*:

$\text{inext } n (I \otimes k) = (\text{if } n \text{ mod } k = 0 \text{ then } (\text{inext } (n \text{ div } k) I) * k \text{ else } n)$
 ⟨proof⟩

lemma *iT-Mult-iprev2-if*:

$\text{iprev } n (I \otimes k) = (\text{if } n \text{ mod } k = 0 \text{ then } (\text{iprev } (n \text{ div } k) I) * k \text{ else } n)$
 ⟨proof⟩

corollary *iT-Mult-inext2*:

$n \text{ mod } k = 0 \implies \text{inext } n (I \otimes k) = (\text{inext } (n \text{ div } k) I) * k$
 ⟨proof⟩

corollary *iT-Mult-iprev2*:

$n \text{ mod } k = 0 \implies \text{iprev } n (I \otimes k) = (\text{iprev } (n \text{ div } k) I) * k$
 ⟨proof⟩

lemma *iT-Plus-neg-inext*:

$k \leq n \implies \text{inext } (n - k) (I \oplus - k) = \text{inext } n I - k$
 ⟨proof⟩

thm *subst[OF inext-cut-ge-conv]*

⟨proof⟩

thm *inext-image*

⟨proof⟩

thm *strict-mono-on-subset[OF - Int-lower2]*

⟨proof⟩

thm *sub-left-strict-mono-on*

⟨proof⟩

lemma *iT-Plus-neg-iprev:*

$$\text{iprev } (n - k) (I \oplus - k) = \text{iprev } n (I \downarrow \geq k) - k$$

<proof>

thm *iT-Plus-neg-mem-iff2[THEN iffD2]*

<proof>

thm *strict-mono-on-subset[OF - Int-lower2]*

<proof>

thm *sub-left-strict-mono-on*

<proof>

corollary *iT-Plus-neg-inext2: inext n (I \oplus - k) = inext (n + k) I - k*

<proof>

corollary *iT-Plus-neg-iprev2: iprev n (I \oplus - k) = iprev (n + k) (I \downarrow \geq k) - k*

<proof>

lemma *iT-Minus-inext:*

$$\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{inext } (k - n) (k \ominus I) = k - \text{iprev } n I$$

<proof>

thm *iT-Minus-imirror-conv*

<proof>

thm *iT-Plus-inext*

<proof>

thm *iprev-mono[THEN order-trans, of n iMin (I \downarrow \leq k) + Max (I \downarrow \leq k) I]*

<proof>

corollary *iT-Minus-inext2:*

$$\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{inext } n (k \ominus I) = k - \text{iprev } (k - n) I$$

<proof>

lemma *iT-Minus-iprev:*

$$\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{iprev } (k - n) (k \ominus I) = k - \text{inext } n (I \downarrow \leq k)$$

<proof>

lemma *iT-Minus-iprev2:*

$$\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{iprev } n (k \ominus I) = k - \text{inext } (k - n) (I \downarrow \leq k)$$

<proof>

lemma *iT-Plus-inext-nth: I \neq \{\} \implies (I \oplus k) \rightarrow n = (I \rightarrow n) + k*

<proof>

lemma *iT-Plus-iprev-nth: \llbracket finite I; I \neq \{\} \rrbracket \implies (I \oplus k) \leftarrow n = (I \leftarrow n) + k*

<proof>

lemma *iT-Mult-inext-nth: I \neq \{\} \implies (I \otimes k) \rightarrow n = (I \rightarrow n) * k*

<proof>

lemma *iT-Mult-iprev-nth*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies (I \otimes k) \leftarrow n = (I \leftarrow n) * k$
 ⟨proof⟩

lemma *iT-Plus-neg-inext-nth*:

$I \oplus - k \neq \{\} \implies (I \oplus - k) \rightarrow n = (I \downarrow \geq k \rightarrow n) - k$
 ⟨proof⟩

thm *iT-Plus-neg-cut-eq[of k k I]*

⟨proof⟩

lemma *iT-Plus-neg-iprev-nth*:

$\llbracket \text{finite } I; I \oplus - k \neq \{\} \rrbracket \implies (I \oplus - k) \leftarrow n = (I \downarrow \geq k \leftarrow n) - k$
 ⟨proof⟩

lemma *iT-Minus-inext-nth*:

$k \ominus I \neq \{\} \implies (k \ominus I) \rightarrow n = k - ((I \downarrow \leq k) \leftarrow n)$
 ⟨proof⟩

lemma *iT-Minus-iprev-nth*:

$k \ominus I \neq \{\} \implies (k \ominus I) \leftarrow n = k - ((I \downarrow \leq k) \rightarrow n)$
 ⟨proof⟩

lemma *iT-Div-ge-inext-nth*:

$\llbracket I \neq \{\}; \forall x \in I. \forall y \in I. x < y \longrightarrow x + k \leq y \rrbracket \implies$
 $(I \circ k) \rightarrow n = (I \rightarrow n) \text{ div } k$
 ⟨proof⟩

lemma *iT-Div-mod-inext-nth*:

$\llbracket I \neq \{\}; \forall x \in I. \forall y \in I. x \bmod k = y \bmod k \rrbracket \implies$
 $(I \circ k) \rightarrow n = (I \rightarrow n) \text{ div } k$
 ⟨proof⟩

lemma *iT-Div-ge-iprev-nth*:

$\llbracket \text{finite } I; I \neq \{\}; \forall x \in I. \forall y \in I. x < y \longrightarrow x + k \leq y \rrbracket \implies$
 $(I \circ k) \leftarrow n = (I \leftarrow n) \text{ div } k$
 ⟨proof⟩

lemma *iT-Div-mod-iprev-nth*:

$\llbracket \text{finite } I; I \neq \{\}; \forall x \in I. \forall y \in I. x \bmod k = y \bmod k \rrbracket \implies$
 $(I \circ k) \leftarrow n = (I \leftarrow n) \text{ div } k$
 ⟨proof⟩

2.4 Cardinality of intervals with interval operators

lemma *iT-Plus-card*: $\text{card } (I \oplus k) = \text{card } I$

⟨proof⟩

lemma *iT-Mult-card*: $0 < k \implies \text{card } (I \otimes k) = \text{card } I$

⟨proof⟩

lemma *iT-Plus-neg-card*: $\text{card } (I \oplus - k) = \text{card } (I \downarrow \geq k)$

⟨proof⟩

thm *Int-lower2*

thm *subset-inj-on[OF - Int-lower2]*

⟨proof⟩

thm *sub-left-inj-on*

⟨proof⟩

lemma *iT-Plus-neg-card-le*: $\text{card } (I \oplus - k) \leq \text{card } I$

⟨proof⟩

thm *nat-cut-ge-finite-iff card-infinite*

⟨proof⟩

lemma *iT-Minus-card*: $\text{card } (k \ominus I) = \text{card } (I \downarrow \leq k)$

⟨proof⟩

thm *subset-inj-on[OF - Int-lower2]*

⟨proof⟩

lemma *iT-Minus-card-le*: $\text{finite } I \implies \text{card } (k \ominus I) \leq \text{card } I$

⟨proof⟩

lemma *iT-Div-0-card-if*:

$\text{card } (I \oslash 0) = (\text{if } I = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$

⟨proof⟩

lemma *Int-empty-setsum*:

$(\sum k \leq (n::\text{nat}). \text{if } \{\} \cap (I k) = \{\} \text{ then } 0 \text{ else } \text{Suc } 0) = 0$

⟨proof⟩

lemma *iT-Div-mod-partition-card*:

$\text{card } (I \cap [n * d.., d - \text{Suc } 0] \oslash d) =$

$(\text{if } I \cap [n * d.., d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$

⟨proof⟩

thm *iT-Div-mem-iff*

⟨proof⟩

thm *le-less-imp-div*

⟨proof⟩

thm *iT-Div-mem-iff-Int*

lemma *iT-Div-conv-count*:

$0 < d \implies I \oslash d = \{k. I \cap [k * d.., d - \text{Suc } 0] \neq \{\}\}$

⟨proof⟩

thm *iT-Div-mem-iff-Int*

⟨proof⟩

lemma *iT-Div-conv-count2*:

$\llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$

$I \oslash d = \{k. k \leq n \wedge I \cap [k * d.., d - \text{Suc } 0] \neq \{\}\}$

⟨proof⟩

thm *div-less-conv*

⟨proof⟩

lemma *mod-partition-count-Suc*:

$\{k. k \leq \text{Suc } n \wedge I \cap [k * d.., d - \text{Suc } 0] \neq \{\}\} =$

$\{k. k \leq n \wedge I \cap [k * d.., d - \text{Suc } 0] \neq \{\}\} \cup$

$(\text{if } I \cap [\text{Suc } n * d.., d - \text{Suc } 0] \neq \{\} \text{ then } \{\text{Suc } n\} \text{ else } \{\})$

⟨proof⟩

lemma *iT-Div-card*:

$\bigwedge I. \llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $\text{card } (I \otimes d) = (\sum k \leq n.$
 $\text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$

<proof>

thm *iT-Div-conv-count2*

<proof>

thm *mod-partition-count-Suc*

<proof>

find-theorems - *div - <= - name: conv*

<proof>

thm *div-le-conv*

<proof>

thm *cut-le-Int-conv*

<proof>

thm *iT-Div-card*

thm *iT-Div-card[OF - - le-reft, of d i]*

corollary *iT-Div-card-Suc*:

$\bigwedge I. \llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $\text{card } (I \otimes d) = (\sum k < \text{Suc } n.$
 $\text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$

<proof>

corollary *iT-Div-Max-card*: $\llbracket 0 < d; \text{finite } I \rrbracket \implies$

$\text{card } (I \otimes d) = (\sum k \leq \text{Max } I \text{ div } d.$
 $\text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$

<proof>

lemma *iT-Div-card-le*: $0 < k \implies \text{card } (I \otimes k) \leq \text{card } I$

<proof>

thm *iT-Div-def*

lemma *iT-Div-card-inj-on*:

inj-on $(\lambda n. n \text{ div } k) I \implies \text{card } (I \otimes k) = \text{card } I$

<proof>

thm *mod-Suc*

lemma *mod-Suc'*:

$0 < n \implies \text{Suc } m \text{ mod } n = (\text{if } m \text{ mod } n < n - \text{Suc } 0 \text{ then } \text{Suc } (m \text{ mod } n) \text{ else } 0)$

<proof>

lemma *div-Suc*:

$0 < n \implies \text{Suc } m \text{ div } n = (\text{if } \text{Suc } (m \text{ mod } n) = n \text{ then } \text{Suc } (m \text{ div } n) \text{ else } m \text{ div } n)$

n)

<proof>

thm *le-neq-trans*[*OF mod-less-divisor*[*THEN Suc-leI*]]

<proof>

lemma *div-Suc'*:

$0 < n \implies \text{Suc } m \text{ div } n = (\text{if } m \text{ mod } n < n - \text{Suc } 0 \text{ then } m \text{ div } n \text{ else } \text{Suc } (m \text{ div } n))$

<proof>

lemma *iT-Div-card-ge-aux*:

$\bigwedge I. \llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$

$\text{card } I \text{ div } d + (\text{if } \text{card } I \text{ mod } d = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq \text{card } (I \oslash d)$

<proof>

thm *div-le-conv*[*THEN iffD1*]

<proof>

thm *nat-card-le-Max*

thm *order-trans*[*OF nat-card-le-Max*]

<proof>

thm *div-le-mono*[*of - - d*]

<proof>

thm *subst*[**where** $t=I$ **and** $s=I \cap [\dots n * d + d - \text{Suc } 0] \cup I \cap [\text{Suc } n * d \dots, d - \text{Suc } 0]$]

<proof>

thm *subset-atMost-Max-le-conv*

thm *le-less-div-conv*

<proof>

thm *card-Un-disjoint*

<proof>

thm *iT-Div-mod-partition-card*

<proof>

thm *Int-card2*[*OF iIN-finite*]

thm *order-trans*[*OF Int-card2*[*OF iIN-finite*]]

<proof>

thm *Int-card2*[*OF iIN-finite, rule-format*]

<proof>

thm *Max-Int-le2*[*OF - iTILL-finite*]

thm *order-trans*[*OF Max-Int-le2*[*OF - iTILL-finite*]]

<proof>

thm *card-Un-disjoint*

<proof>

thm *mod-partition-iT-Div-Int-one-segment*

<proof>

thm *mod-partition-iT-Div-Int-one-segment*

<proof>

thm *mod-partition-iT-Div-Int2*

<proof>

thm *disjoint-iff-in-not-in1*

<proof>

thm *iIN-div-mod-eq-0*

⟨proof⟩
thm *iT-Div-mod-partition-card*
 ⟨proof⟩
thm *add-le-divisor-imp-le-Suc-div*
 ⟨proof⟩
thm *Int-card2*[OF *iIN-finite*, THEN *le-trans*]
 ⟨proof⟩
thm *order-trans*[where $x = \text{card } I \text{ div } d + (\text{if } \text{card } I \text{ mod } d \neq 0 \text{ then } \text{Suc } 0 \text{ else } 0)$]
 ⟨proof⟩
thm *div-add1-eq1-mod-0-left*
 ⟨proof⟩
thm *add-le-divisor-imp-le-Suc-div*
 ⟨proof⟩

lemma *iT-Div-card-ge*:
 $\text{card } I \text{ div } d + (\text{if } \text{card } I \text{ mod } d = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq \text{card } (I \otimes d)$
 ⟨proof⟩
thm *iT-Div-card-ge-aux*[OF - - *order-refl*]
 ⟨proof⟩
corollary *iT-Div-card-ge-div*: $\text{card } I \text{ div } d \leq \text{card } (I \otimes d)$
 ⟨proof⟩

There is no better lower bound function f for $i \otimes d$ with $\text{card } i$ and d as arguments.

lemma *iT-Div-card-ge--is-maximal-lower-bound*:
 $\forall I d. \text{card } I \text{ div } d + (\text{if } \text{card } I \text{ mod } d = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq f (\text{card } I) d \wedge$
 $f (\text{card } I) d \leq \text{card } (I \otimes d) \implies$
 $f (\text{card } (I::\text{nat set})) d = \text{card } I \text{ div } d + (\text{if } \text{card } I \text{ mod } d = 0 \text{ then } 0 \text{ else } \text{Suc } 0)$
 ⟨proof⟩
thm *iTILL-div iTILL-card*
 ⟨proof⟩
thm *div-diff1-eq1*[of *Suc 0 d card I*]
 ⟨proof⟩
thm *iT-Div-card-ge--is-maximal-lower-bound*[rule-format]

thm *iT-Plus-card*
lemma *iT-Plus-icard*: $\text{icard } (I \oplus k) = \text{icard } I$
 ⟨proof⟩

thm *iT-Mult-card*
lemma *iT-Mult-icard*: $0 < k \implies \text{icard } (I \otimes k) = \text{icard } I$
 ⟨proof⟩

thm *iT-Plus-neg-card*
lemma *iT-Plus-neg-icard*: $\text{icard } (I \oplus - k) = \text{icard } (I \downarrow \geq k)$
 ⟨proof⟩
thm *iT-Plus-neg-finite-iff cut-ge-finite*

<proof>

thm *iT-Plus-neg-card-le*

lemma *iT-Plus-neg-icard-le*: $\text{icard } (I \oplus - k) \leq \text{icard } I$

<proof>

thm *iT-Minus-card*

lemma *iT-Minus-icard*: $\text{icard } (k \ominus I) = \text{icard } (I \downarrow \leq k)$

<proof>

thm *iT-Minus-card-le*

lemma *iT-Minus-icard-le*: $\text{icard } (k \ominus I) \leq \text{icard } I$

<proof>

thm *iT-Div-0-card-if*

lemma *iT-Div-0-icard-if*: $\text{icard } (I \oslash 0) = \text{Fin } (\text{if } I = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$

thm *iT-Div-0-finite*

<proof>

thm *iT-Div-mod-partition-card*

lemma *iT-Div-mod-partition-icard*:

$\text{icard } (I \cap [n * d \dots, d - \text{Suc } 0] \oslash d) =$

$\text{Fin } (\text{if } I \cap [n * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$

<proof>

thm *iT-Div-card*

lemma *iT-Div-icard*:

$\llbracket 0 < d; \text{finite } I \implies \text{Max } I \text{ div } d \leq n \rrbracket \implies$

$\text{icard } (I \oslash d) =$

$(\text{if } \text{finite } I \text{ then } \text{Fin } (\sum k \leq n. \text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0) \text{ else } \infty)$

<proof>

thm *iT-Div-Max-card*

corollary *iT-Div-Max-icard*: $0 < d \implies$

$\text{icard } (I \oslash d) = (\text{if } \text{finite } I$

$\text{then } \text{Fin } (\sum k \leq \text{Max } I \text{ div } d. \text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0) \text{ else } \infty)$

<proof>

thm *iT-Div-card-le*

lemma *iT-Div-icard-le*: $0 < k \implies \text{icard } (I \oslash k) \leq \text{icard } I$

<proof>

thm *iT-Div-card-inj-on*

lemma *iT-Div-icard-inj-on*: $\text{inj-on } (\lambda n. n \text{ div } k) I \implies \text{icard } (I \oslash k) = \text{icard } I$

<proof>

thm *iT-Div-card-ge*

lemma *iT-Div-icard-ge*: $\text{icard } I \text{ div } (\text{Fin } d) + \text{Fin } (\text{if } \text{icard } I \text{ mod } (\text{Fin } d) = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq \text{icard } (I \odot d)$

<proof>

thm *iT-Div-card-ge-div*

corollary *iT-Div-icard-ge-div*: $\text{icard } I \text{ div } (\text{Fin } d) \leq \text{icard } (I \odot d)$

<proof>

thm *iT-Div-card-ge--is-maximal-lower-bound*

lemma *iT-Div-icard-ge--is-maximal-lower-bound*:

$\forall I d. \text{icard } I \text{ div } (\text{Fin } d) + \text{Fin } (\text{if } \text{icard } I \text{ mod } (\text{Fin } d) = 0 \text{ then } 0 \text{ else } \text{Suc } 0)$

$\leq f (\text{icard } I) d \wedge$

$f (\text{icard } I) d \leq \text{icard } (I \odot d) \implies$

$f (\text{icard } (I::\text{nat set})) d =$

$\text{icard } I \text{ div } (\text{Fin } d) + \text{Fin } (\text{if } \text{icard } I \text{ mod } (\text{Fin } d) = 0 \text{ then } 0 \text{ else } \text{Suc } 0)$

<proof>

thm *iT-Div-card-ge--is-maximal-lower-bound*

<proof>

thm *iT-Div-card-ge--is-maximal-lower-bound*

<proof>

2.5 Results about sets of intervals

2.5.1 Set of intervals without and with empty interval

definition *iFROM-UN-set* :: $(\text{nat set}) \text{ set where}$

$iFROM\text{-UN}\text{-set} \equiv \bigcup n. \{[n..]\}$

definition *iTILL-UN-set* :: $(\text{nat set}) \text{ set where}$

$iTILL\text{-UN}\text{-set} \equiv \bigcup n. \{[..n]\}$

definition *iIN-UN-set* :: $(\text{nat set}) \text{ set where}$

$iIN\text{-UN}\text{-set} \equiv \bigcup n d. \{[n..,d]\}$

definition *iMOD-UN-set* :: $(\text{nat set}) \text{ set where}$

$iMOD\text{-UN}\text{-set} \equiv \bigcup r m. \{[r, \text{mod } m]\}$

definition *iMODb-UN-set* :: $(\text{nat set}) \text{ set where}$

$iMODb\text{-UN}\text{-set} \equiv \bigcup r m c. \{[r, \text{mod } m, c]\}$

definition *iFROM-set* :: $(\text{nat set}) \text{ set where}$

$iFROM\text{-set} \equiv \{[n..] \mid n. \text{True}\}$

definition *iTILL-set* :: $(\text{nat set}) \text{ set where}$

$iTILL\text{-set} \equiv \{[..n] \mid n. \text{True}\}$

definition *iIN-set* :: $(\text{nat set}) \text{ set where}$

$iIN\text{-set} \equiv \{[n..,d] \mid n d. \text{True}\}$

definition *iMOD-set* :: $(\text{nat set}) \text{ set where}$

$iMOD\text{-set} \equiv \{[r, \text{mod } m] \mid r m. \text{True}\}$

definition *iMODb-set* :: $(\text{nat set}) \text{ set where}$

$iMODb\text{-set} \equiv \{[r, \text{mod } m, c] \mid r m c. \text{True}\}$

definition *iMOD2-set* :: $(\text{nat set}) \text{ set where}$

$iMOD2\text{-set} \equiv \{[r, \text{mod } m] \mid r m. 2 \leq m\}$

definition $iMODb2\text{-set} :: (\text{nat set}) \text{ set}$ **where**
 $iMODb2\text{-set} \equiv \{[r, \text{mod } m, c] \mid r \text{ mod } m \text{ c. } 2 \leq m \wedge 1 \leq c\}$

definition $iMOD2\text{-UN-set} :: (\text{nat set}) \text{ set}$ **where**
 $iMOD2\text{-UN-set} \equiv \bigcup r. \bigcup m \in \{2..\}. \{[r, \text{mod } m]\}$

definition $iMODb2\text{-UN-set} :: (\text{nat set}) \text{ set}$ **where**
 $iMODb2\text{-UN-set} \equiv \bigcup r. \bigcup m \in \{2..\}. \bigcup c \in \{1..\}. \{[r, \text{mod } m, c]\}$

lemmas $i\text{-set-defs} =$
 $iFROM\text{-set-def } iTILL\text{-set-def } iIN\text{-set-def}$
 $iMOD\text{-set-def } iMODb\text{-set-def}$
 $iMOD2\text{-set-def } iMODb2\text{-set-def}$

lemmas $i\text{-UN-set-defs} =$
 $iFROM\text{-UN-set-def } iTILL\text{-UN-set-def } iIN\text{-UN-set-def}$
 $iMOD\text{-UN-set-def } iMODb\text{-UN-set-def}$
 $iMOD2\text{-UN-set-def } iMODb2\text{-UN-set-def}$

lemma $iFROM\text{-set-UN-set-eq}: iFROM\text{-set} = iFROM\text{-UN-set}$
 $\langle \text{proof} \rangle$

lemma
 $iTILL\text{-set-UN-set-eq}: iTILL\text{-set} = iTILL\text{-UN-set}$ **and**
 $iIN\text{-set-UN-set-eq}: iIN\text{-set} = iIN\text{-UN-set}$ **and**
 $iMOD\text{-set-UN-set-eq}: iMOD\text{-set} = iMOD\text{-UN-set}$ **and**
 $iMODb\text{-set-UN-set-eq}: iMODb\text{-set} = iMODb\text{-UN-set}$
 $\langle \text{proof} \rangle$

lemma $iMOD2\text{-set-UN-set-eq}: iMOD2\text{-set} = iMOD2\text{-UN-set}$
 $\langle \text{proof} \rangle$

lemma $iMODb2\text{-set-UN-set-eq}: iMODb2\text{-set} = iMODb2\text{-UN-set}$
 $\langle \text{proof} \rangle$

lemmas $i\text{-set-i-UN-set-sets-eq} =$
 $iFROM\text{-set-UN-set-eq}$
 $iTILL\text{-set-UN-set-eq}$
 $iIN\text{-set-UN-set-eq}$
 $iMOD\text{-set-UN-set-eq}$
 $iMODb\text{-set-UN-set-eq}$
 $iMOD2\text{-set-UN-set-eq}$
 $iMODb2\text{-set-UN-set-eq}$

thm $i\text{-set-i-UN-set-sets-eq}$

lemma $iMOD2\text{-set-iMOD-set-subset}: iMOD2\text{-set} \subseteq iMOD\text{-set}$
 $\langle \text{proof} \rangle$

lemma $iMODb2\text{-set-iMODb-set-subset}: iMODb2\text{-set} \subseteq iMODb\text{-set}$
 $\langle \text{proof} \rangle$

definition $i\text{-set} :: (\text{nat set}) \text{ set}$ **where**

$i\text{-set} \equiv iFROM\text{-set} \cup iTILL\text{-set} \cup iIN\text{-set} \cup iMOD\text{-set} \cup iMODb\text{-set}$

definition $i\text{-UN-set} :: (\text{nat set}) \text{ set}$ **where**

$i\text{-UN-set} \equiv i\text{FROM-UN-set} \cup i\text{TILL-UN-set} \cup i\text{IN-UN-set} \cup i\text{MOD-UN-set} \cup i\text{MODb-UN-set}$

Minimal definitions for $i\text{-set}$ and $i\text{-set}$

definition $i\text{-set-min} :: (\text{nat set}) \text{ set where}$

$i\text{-set-min} \equiv i\text{FROM-set} \cup i\text{IN-set} \cup i\text{MOD2-set} \cup i\text{MODb2-set}$

definition $i\text{-UN-set-min} :: (\text{nat set}) \text{ set where}$

$i\text{-UN-set-min} \equiv i\text{FROM-UN-set} \cup i\text{IN-UN-set} \cup i\text{MOD2-UN-set} \cup i\text{MODb2-UN-set}$

definition $i\text{-set0} :: (\text{nat set}) \text{ set where}$

$i\text{-set0} \equiv \text{insert } \{\} \ i\text{-set}$

thm $i\text{-set0-def}$

thm $i\text{-set-i-UN-set-sets-eq}$

lemma $i\text{-set-i-UN-set-eq}: i\text{-set} = i\text{-UN-set}$

$\langle \text{proof} \rangle$

lemma $i\text{-set-min-i-UN-set-min-eq}: i\text{-set-min} = i\text{-UN-set-min}$

$\langle \text{proof} \rangle$

lemma $i\text{-set-min-eq}: i\text{-set} = i\text{-set-min}$

$\langle \text{proof} \rangle$

thm $Un\text{-mono}$

$\langle \text{proof} \rangle$

corollary $i\text{-UN-set-i-UN-min-set-eq}: i\text{-UN-set} = i\text{-UN-set-min}$

$\langle \text{proof} \rangle$

thm

$i\text{-set-def } i\text{-set-min-def}$

$i\text{-set-min-eq}$

lemma $i\text{-set-min-is-minimal-let}:$

$\text{let } s1 = i\text{FROM-set}; s2 = i\text{IN-set}; s3 = i\text{MOD2-set}; s4 = i\text{MODb2-set} \text{ in}$

$s1 \cap s2 = \{\} \wedge s1 \cap s3 = \{\} \wedge s1 \cap s4 = \{\} \wedge$

$s2 \cap s3 = \{\} \wedge s2 \cap s4 = \{\} \wedge s3 \cap s4 = \{\}$

$\langle \text{proof} \rangle$

lemmas $i\text{-set-min-is-minimal} = i\text{-set-min-is-minimal-let}[\text{simplified}]$

thm $i\text{-set-min-is-minimal}$

thm conjunct1

thm

$i\text{-set-min-is-minimal}[\text{THEN } \text{conjunct1}]$

$i\text{-set-min-is-minimal}[\text{THEN } \text{conjunct1}[\text{OF } \text{conjunct2}]]$

$i\text{-set-min-is-minimal}[\text{THEN } \text{conjunct1}[\text{OF } \text{conjunct2}[\text{OF } \text{conjunct2}]]]$

$i\text{-set-min-is-minimal}[\text{THEN } \text{conjunct1}[\text{OF } \text{conjunct2}[\text{OF } \text{conjunct2}[\text{OF } \text{conjunct2}]]]]$

$i\text{-set-min-is-minimal}[\text{THEN } \text{conjunct1}[\text{OF } \text{conjunct2}[\text{OF } \text{conjunct2}[\text{OF } \text{conjunct2}[\text{OF } \text{conjunct2}]]]]]$

i-set-min-is-minimal[*THEN conjunct2*[*OF conjunct2*[*OF conjunct2*[*OF conjunct2*[*OF conjunct2*]]]]]]

thm

i-set-def

i-set-defs

inductive-set

i-set-ind:: (nat set) set

where

i-set-ind-FROM[*intro!*]: [*n...*] ∈ *i-set-ind*

| *i-set-ind-TILL*[*intro!*]: [*..n*] ∈ *i-set-ind*

| *i-set-ind-IN*[*intro!*]: [*n...d*] ∈ *i-set-ind*

| *i-set-ind-MOD*[*intro!*]: [*r, mod m*] ∈ *i-set-ind*

| *i-set-ind-MODb*[*intro!*]: [*r, mod m, c*] ∈ *i-set-ind*

inductive-set

i-set0-ind :: (nat set) set

where

i-set0-ind-empty[*intro!*]: {} ∈ *i-set0-ind*

| *i-set0-ind-i-set*[*intro!*]: *I* ∈ *i-set-ind* ⇒ *I* ∈ *i-set0-ind*

The introduction rule *i-set0-ind-i-set* is not declared a safe introduction rule, because it would disturb the correct usage of the *safe* method.

lemma *i-set-ind-subset-i-set0-ind*: *i-set-ind* ⊆ *i-set0-ind*

<proof>

thm

i-set-ind-FROM

i-set-ind-TILL

i-set-ind-IN

i-set-ind-MOD

i-set-ind-MODb

lemma

i-set0-ind-FROM[*intro!*]: [*n...*] ∈ *i-set0-ind* **and**

i-set0-ind-TILL[*intro!*]: [*..n*] ∈ *i-set0-ind* **and**

i-set0-ind-IN[*intro!*]: [*n...d*] ∈ *i-set0-ind* **and**

i-set0-ind-MOD[*intro!*]: [*r, mod m*] ∈ *i-set0-ind* **and**

i-set0-ind-MODb[*intro!*]: [*r, mod m, c*] ∈ *i-set0-ind*

<proof>

lemmas *i-set0-ind-intros2* =

i-set0-ind-empty

i-set0-ind-FROM

i-set0-ind-TILL

i-set0-ind-IN

i-set0-ind-MOD

i-set0-ind-MODb

thm *i-set0-ind-intros2*

lemma *i-set-i-set-ind-eq*: $i\text{-set} = i\text{-set-ind}$
 ⟨proof⟩

lemma *i-set0-i-set0-ind-eq*: $i\text{-set0} = i\text{-set0-ind}$
 ⟨proof⟩

thm *i-set-i-set-ind-eq*
 ⟨proof⟩

lemma *i-set-imp-not-empty*: $I \in i\text{-set} \implies I \neq \{\}$
 ⟨proof⟩

lemma *i-set0-i-set-mem-conv*: $(I \in i\text{-set0}) = (I \in i\text{-set} \vee I = \{\})$
 ⟨proof⟩

lemma *i-set-i-set0-mem-conv*: $(I \in i\text{-set}) = (I \in i\text{-set0} \wedge I \neq \{\})$
 ⟨proof⟩

lemma *i-set0-i-set-conv*: $i\text{-set0} - \{\{\}\} = i\text{-set}$
 ⟨proof⟩

corollary *i-set-subset-i-set0*: $i\text{-set} \subseteq i\text{-set0}$
 ⟨proof⟩

lemma *i-set-singleton*: $\{a\} \in i\text{-set}$
 ⟨proof⟩

lemma *i-set0-singleton*: $\{a\} \in i\text{-set0}$
 ⟨proof⟩

thm *i-set-ind.intros*

corollary

i-set-FROM[intro!]: $[n\dots] \in i\text{-set}$ **and**
i-set-TILL[intro!]: $[\dots n] \in i\text{-set}$ **and**
i-set-IN[intro!]: $[n\dots, d] \in i\text{-set}$ **and**
i-set-MOD[intro!]: $[r, \text{mod } m] \in i\text{-set}$ **and**
i-set-MODb[intro!]: $[r, \text{mod } m, c] \in i\text{-set}$

thm *i-set-i-set-ind-eq*
 ⟨proof⟩

lemmas *i-set-intros* =

i-set-FROM

i-set-TILL

i-set-IN

i-set-MOD

i-set-MODb

thm *i-set0-ind-intros2*

lemma

i-set0-empty[intro!]: $\{\} \in i\text{-set0}$ **and**
i-set0-FROM[intro!]: $[n\dots] \in i\text{-set0}$ **and**
i-set0-TILL[intro!]: $[\dots n] \in i\text{-set0}$ **and**
i-set0-IN[intro!]: $[n\dots, d] \in i\text{-set0}$ **and**
i-set0-MOD[intro!]: $[r, \text{mod } m] \in i\text{-set0}$ **and**

$i\text{-set0-MODb}[\text{intro!}] : [r, \text{mod } m, c] \in i\text{-set0}$
 ⟨proof⟩

lemmas $i\text{-set0-intros} =$

$i\text{-set0-empty}$

$i\text{-set0-FROM}$

$i\text{-set0-TILL}$

$i\text{-set0-IN}$

$i\text{-set0-MOD}$

$i\text{-set0-MODb}$

thm $i\text{-set0-intros}$

lemma $i\text{-set-infinite-as-iMOD}$:

$\llbracket I \in i\text{-set}; \text{infinite } I \rrbracket \implies \exists r m. I = [r, \text{mod } m]$
 ⟨proof⟩

lemma $i\text{-set-finite-as-iMODb}$:

$\llbracket I \in i\text{-set}; \text{finite } I \rrbracket \implies \exists r m c. I = [r, \text{mod } m, c]$
 ⟨proof⟩

lemma $i\text{-set-as-iMOD-iMODb}$:

$I \in i\text{-set} \implies (\exists r m. I = [r, \text{mod } m]) \vee (\exists r m c. I = [r, \text{mod } m, c])$
 ⟨proof⟩

2.5.2 Interval sets are closed under cutting

lemma $i\text{-set-cut-le-ge-closed-disj}$:

$\llbracket I \in i\text{-set}; t \in I; \text{cut-op} = \text{op } \downarrow \leq \vee \text{cut-op} = \text{op } \downarrow \geq \rrbracket \implies$
 $\text{cut-op } I t \in i\text{-set}$
 ⟨proof⟩

thm $i\text{-set-cut-le-ge-closed-disj}[\text{of } - - \text{op } \downarrow \geq, \text{ simplified}]$

corollary

$i\text{-set-cut-le-closed}$: $\llbracket I \in i\text{-set}; t \in I \rrbracket \implies I \downarrow \leq t \in i\text{-set}$ **and**

$i\text{-set-cut-ge-closed}$: $\llbracket I \in i\text{-set}; t \in I \rrbracket \implies I \downarrow \geq t \in i\text{-set}$

⟨proof⟩

lemmas $i\text{-set-cut-le-ge-closed} = i\text{-set-cut-le-closed } i\text{-set-cut-ge-closed}$

thm $i\text{-set-cut-le-ge-closed}$

lemma $i\text{-set0-cut-closed-disj}$:

$\llbracket I \in i\text{-set0};$
 $\text{cut-op} = \text{op } \downarrow \leq \vee \text{cut-op} = \text{op } \downarrow \geq \vee$
 $\text{cut-op} = \text{op } \downarrow < \vee \text{cut-op} = \text{op } \downarrow > \rrbracket \implies$
 $\text{cut-op } I t \in i\text{-set0}$
 ⟨proof⟩

thm $i\text{-set0-ind-empty}$

thm *ssubst*[*OF set-restriction-empty*, **where** $P = \lambda x. x \in i\text{-set0-ind}$]
 ⟨*proof*⟩
thm *i-cut-set-restriction-disj*
 ⟨*proof*⟩

thm *i-set0-cut-closed-disj*[*of - op* \downarrow >, *simplified*]

corollary

i-set0-cut-le-closed: $I \in i\text{-set0} \implies I \downarrow \leq t \in i\text{-set0}$ **and**
i-set0-cut-less-closed: $I \in i\text{-set0} \implies I \downarrow < t \in i\text{-set0}$ **and**
i-set0-cut-ge-closed: $I \in i\text{-set0} \implies I \downarrow \geq t \in i\text{-set0}$ **and**
i-set0-cut-greater-closed: $I \in i\text{-set0} \implies I \downarrow > t \in i\text{-set0}$
 ⟨*proof*⟩

thm *i-set0-ind.intros*

thm

i-set0-FROM[*THEN* *i-set0-cut-closed-disj*[*of - op* \downarrow <, *simplified*]]
i-set0-TILL[*THEN* *i-set0-cut-closed-disj*[*of - op* \downarrow >, *simplified*]]
i-set0-IN[*THEN* *i-set0-cut-closed-disj*[*of - op* \downarrow <=, *simplified*]]
i-set0-MOD[*THEN* *i-set0-cut-closed-disj*[*of - op* \downarrow >, *simplified*]]
i-set0-MODb[*THEN* *i-set0-cut-closed-disj*[*of - op* \downarrow >, *simplified*]]

lemmas *i-set0-cut-closed* =

i-set0-cut-le-closed
i-set0-cut-less-closed
i-set0-cut-ge-closed
i-set0-cut-greater-closed

thm *i-set0-cut-closed*

2.5.3 Interval sets are closed under addition and multiplication

lemma *i-set-Plus-closed*: $I \in i\text{-set} \implies I \oplus k \in i\text{-set}$
 ⟨*proof*⟩

lemma *i-set-Mult-closed*: $I \in i\text{-set} \implies I \otimes k \in i\text{-set}$
 ⟨*proof*⟩

lemma *i-set0-Plus-closed*: $I \in i\text{-set0} \implies I \oplus k \in i\text{-set0}$
 ⟨*proof*⟩

lemma *i-set0-Mult-closed*: $I \in i\text{-set0} \implies I \otimes k \in i\text{-set0}$
 ⟨*proof*⟩

2.5.4 Interval sets are closed with certain conditions under subtraction

lemma *i-set-Plus-neg-closed*:

[$I \in i\text{-set}; \exists x \in I. k \leq x$] $\implies I \oplus -k \in i\text{-set}$
 ⟨*proof*⟩

lemma *i-set-Minus-closed*:

$\llbracket I \in i\text{-set}; i\text{Min } I \leq k \rrbracket \implies k \ominus I \in i\text{-set}$
 ⟨proof⟩

lemma *i-set0-Plus-neg-closed*: $I \in i\text{-set0} \implies I \oplus - k \in i\text{-set0}$
 ⟨proof⟩

lemma *i-set0-Minus-closed*: $I \in i\text{-set0} \implies k \ominus I \in i\text{-set0}$
 ⟨proof⟩

lemmas *i-set-IntOp-closed* =

i-set-Plus-closed
i-set-Mult-closed
i-set-Plus-neg-closed
i-set-Minus-closed

lemmas *i-set0-IntOp-closed* =

i-set0-Plus-closed
i-set0-Mult-closed
i-set0-Plus-neg-closed
i-set0-Minus-closed

thm *i-set-IntOp-closed*

thm *i-set0-IntOp-closed*

2.5.5 Interval sets are not closed under division

thm *i-set-as-iMOD-iMODb*

thm *iMOD-div-mod-gr0-not-ex*

lemma *iMOD-div-mod-gr0-not-in-i-set*:

$\llbracket 0 < k; k < m; 0 < m \bmod k \rrbracket \implies [r, \text{mod } m] \otimes k \notin i\text{-set}$
 ⟨proof⟩

thm *i-set-infinite-as-iMOD*

⟨proof⟩

lemma *iMODb-div-mod-gr0-not-in-i-set*:

$\llbracket 0 < k; k < m; 0 < m \bmod k; k \leq c \rrbracket \implies [r, \text{mod } m, c] \otimes k \notin i\text{-set}$
 ⟨proof⟩

lemma $[0, \text{mod } 3] \otimes 2 \notin i\text{-set}$

⟨proof⟩

lemma *i-set-Div-not-closed*: $\text{Suc } 0 < k \implies \exists I \in i\text{-set}. I \otimes k \notin i\text{-set}$

⟨proof⟩

lemma *i-set0-Div-not-closed*: $\text{Suc } 0 < k \implies \exists I \in i\text{-set0}. I \otimes k \notin i\text{-set0}$

⟨proof⟩

thm

i-set-Div-not-closed

i-set0-Div-not-closed

2.5.6 Sets of intervals closed under division

inductive-set

NatMultiples :: *nat set* \Rightarrow *nat set*

for *F* :: *nat set*

where

NatMultiples-Factor:

$k \in F \Longrightarrow k \in \text{NatMultiples } F$

| *NatMultiples-Product*:

$\llbracket k \in F; m \in \text{NatMultiples } F \rrbracket \Longrightarrow k * m \in \text{NatMultiples } F$

thm *NatMultiples.induct*

lemma *NatMultiples-ex-divisor*: $m \in \text{NatMultiples } F \Longrightarrow \exists k \in F. m \bmod k = 0$

<proof>

lemma *NatMultiples-product-factor*: $\llbracket a \in F; b \in F \rrbracket \Longrightarrow a * b \in \text{NatMultiples } F$

<proof>

lemma *NatMultiples-product-factor-multiple*:

$\llbracket a \in F; b \in \text{NatMultiples } F \rrbracket \Longrightarrow a * b \in \text{NatMultiples } F$

<proof>

lemma *NatMultiples-product-multiple-factor*:

$\llbracket a \in \text{NatMultiples } F; b \in F \rrbracket \Longrightarrow a * b \in \text{NatMultiples } F$

<proof>

lemma *NatMultiples-product-multiple*:

$\llbracket a \in \text{NatMultiples } F; b \in \text{NatMultiples } F \rrbracket \Longrightarrow a * b \in \text{NatMultiples } F$

<proof>

Subset of *i-set* containing only continuous intervals, i. e., without *iMOD* and *iMODb*.

inductive-set *i-set-cont* :: (*nat set*) *set*

where

i-set-cont-FROM[*intro*]: $[n\dots] \in i\text{-set-cont}$

| *i-set-cont-TILL*[*intro*]: $[\dots n] \in i\text{-set-cont}$

| *i-set-cont-IN*[*intro*]: $[n\dots, d] \in i\text{-set-cont}$

definition *i-set0-cont* :: (*nat set*) *set* where

$i\text{-set0-cont} \equiv \text{insert } \{\} i\text{-set-cont}$

lemma *i-set-cont-subset-i-set*: $i\text{-set-cont} \subseteq i\text{-set}$

<proof>

lemma *i-set0-cont-subset-i-set0*: $i\text{-set0-cont} \subseteq i\text{-set0}$

<proof>

Minimal definition of *i-set-cont*

inductive-set *i-set-cont-min* :: (*nat set*) *set*

where

i-set-cont-FROM[*intro*]: $[n\dots] \in i\text{-set-cont-min}$

| *i-set-cont-IN*[*intro*]: $[n\dots, d] \in i\text{-set-cont-min}$

definition *i-set0-cont-min* :: (nat set) set **where**
i-set0-cont-min \equiv insert {} *i-set-cont-min*

lemma *i-set-cont-min-eq*: *i-set-cont* = *i-set-cont-min*
 ⟨proof⟩

inext and *iprev* with continuous intervals

lemma *i-set-cont-inext*:

$\llbracket I \in i\text{-set-cont}; n \in I; \text{finite } I \implies n < \text{Max } I \rrbracket \implies \text{inext } n \ I = \text{Suc } n$
 ⟨proof⟩

lemma *i-set-cont-iprev*:

$\llbracket I \in i\text{-set-cont}; n \in I; i\text{Min } I < n \rrbracket \implies \text{iprev } n \ I = n - \text{Suc } 0$
 ⟨proof⟩

lemma *i-set-cont-inext-less*:

$\llbracket I \in i\text{-set-cont}; n \in I; n < n0; n0 \in I \rrbracket \implies \text{inext } n \ I = \text{Suc } n$
 ⟨proof⟩

thm *order-less-le-trans[OF - Max-ge]*
 ⟨proof⟩

lemma *i-set-cont-iprev-greater*:

$\llbracket I \in i\text{-set-cont}; n \in I; n0 < n; n0 \in I \rrbracket \implies \text{iprev } n \ I = n - \text{Suc } 0$
 ⟨proof⟩

thm *order-le-less-trans[OF iMin-le, of n0]*
 ⟨proof⟩

Sets of modulo intervals

inductive-set *i-set-mult* :: nat \Rightarrow (nat set) set

for *k* :: nat

where

i-set-mult-FROM[intro!]: $[n..] \in i\text{-set-mult } k$
 | *i-set-mult-TILL*[intro!]: $[\dots n] \in i\text{-set-mult } k$
 | *i-set-mult-IN*[intro!]: $[n\dots, d] \in i\text{-set-mult } k$
 | *i-set-mult-MOD*[intro!]: $[r, \text{mod } m * k] \in i\text{-set-mult } k$
 | *i-set-mult-MODb*[intro!]: $[r, \text{mod } m * k, c] \in i\text{-set-mult } k$

definition *i-set0-mult* :: nat \Rightarrow (nat set) set **where**

i-set0-mult *k* \equiv insert {} (*i-set-mult* *k*)

lemma

i-set0-mult-empty[intro!]: {} $\in i\text{-set0-mult } k$ **and**
i-set0-mult-FROM[intro!]: $[n..] \in i\text{-set0-mult } k$ **and**
i-set0-mult-TILL[intro!]: $[\dots n] \in i\text{-set0-mult } k$ **and**
i-set0-mult-IN[intro!]: $[n\dots, d] \in i\text{-set0-mult } k$ **and**
i-set0-mult-MOD[intro!]: $[r, \text{mod } m * k] \in i\text{-set0-mult } k$ **and**
i-set0-mult-MODb[intro!]: $[r, \text{mod } m * k, c] \in i\text{-set0-mult } k$

⟨proof⟩

lemmas *i-set0-mult-intros* =

i-set0-mult-empty
i-set0-mult-FROM

i-set0-mult-TILL

i-set0-mult-IN

i-set0-mult-MOD

i-set0-mult-MODb

thm *i-set0-mult-intros*

lemma *i-set-mult-subset-i-set0-mult*: $i\text{-set-mult } k \subseteq i\text{-set0-mult } k$
 ⟨*proof*⟩

lemma *i-set-mult-subset-i-set*: $i\text{-set-mult } k \subseteq i\text{-set}$
 ⟨*proof*⟩

thm *subsetD[OF i-set-mult-subset-i-set]*

lemma *i-set0-mult-subset-i-set0*: $i\text{-set0-mult } k \subseteq i\text{-set0}$
 ⟨*proof*⟩

thm *order-trans[OF - i-set-subset-i-set0, OF i-set-mult-subset-i-set]*
 ⟨*proof*⟩

lemma *i-set-mult-0-eq-i-set-cont*: $i\text{-set-mult } 0 = i\text{-set-cont}$
 ⟨*proof*⟩

lemma *i-set0-mult-0-eq-i-set0-cont*: $i\text{-set0-mult } 0 = i\text{-set0-cont}$
 ⟨*proof*⟩

lemma *i-set-mult-1-eq-i-set*: $i\text{-set-mult } (\text{Suc } 0) = i\text{-set}$
 ⟨*proof*⟩

lemma *i-set0-mult-1-eq-i-set0*: $i\text{-set0-mult } (\text{Suc } 0) = i\text{-set0}$
 ⟨*proof*⟩

lemma *i-set-mult-imp-not-empty*: $I \in i\text{-set-mult } k \implies I \neq \{\}$
 ⟨*proof*⟩

lemma *iMOD-in-i-set-mult-imp-divisor-mod-0*:
 $\llbracket m \neq \text{Suc } 0; [r, \text{mod } m] \in i\text{-set-mult } k \rrbracket \implies m \text{ mod } k = 0$
 ⟨*proof*⟩

lemma
divisor-mod-0-imp-iMOD-in-i-set-mult: $m \text{ mod } k = 0 \implies [r, \text{mod } m] \in i\text{-set-mult } k$
and

divisor-mod-0-imp-iMODb-in-i-set-mult: $m \text{ mod } k = 0 \implies [r, \text{mod } m, c] \in i\text{-set-mult } k$
 ⟨*proof*⟩

lemma *iMOD-in-i-set-mult--divisor-mod-0-conv*:
 $m \neq \text{Suc } 0 \implies ([r, \text{mod } m] \in i\text{-set-mult } k) = (m \text{ mod } k = 0)$
 ⟨*proof*⟩

lemma *i-set-mult-neq1-subset-i-set*: $k \neq \text{Suc } 0 \implies i\text{-set-mult } k \subset i\text{-set}$
 ⟨proof⟩

lemma *mod-0-imp-i-set-mult-subset*:
 $a \bmod b = 0 \implies i\text{-set-mult } a \subseteq i\text{-set-mult } b$
 ⟨proof⟩

thm *i-set-mult.cases*
 ⟨proof⟩

lemma *i-set-mult-subset-imp-mod-0*:
 $\llbracket a \neq \text{Suc } 0; (i\text{-set-mult } a \subseteq i\text{-set-mult } b) \rrbracket \implies a \bmod b = 0$
 ⟨proof⟩

thm *iMOD-in-i-set-mult-imp-divisor-mod-0[of - 0 b]*
 ⟨proof⟩

lemma *i-set-mult-subset-conv*:
 $a \neq \text{Suc } 0 \implies (i\text{-set-mult } a \subseteq i\text{-set-mult } b) = (a \bmod b = 0)$
 ⟨proof⟩

lemma *i-set-mult-mod-0-div*:
 $\llbracket I \in i\text{-set-mult } k; k \bmod d = 0 \rrbracket \implies I \oslash d \in i\text{-set-mult } (k \text{ div } d)$
 ⟨proof⟩

thm *iT-Div-0[OF i-set-mult-imp-not-empty]*
 ⟨proof⟩

thm *mod-0-imp-mod-mult-left-0*

thm *mod-0-imp-div-mult-right-eq*

thm *iMOD-div iMODb-div*
 ⟨proof⟩

Intervals from *i-set-mult* k remain in *i-set* after division by d a divisor of k .

corollary *i-set-mult-mod-0-div-i-set*:
 $\llbracket I \in i\text{-set-mult } k; k \bmod d = 0 \rrbracket \implies I \oslash d \in i\text{-set}$

thm *subsetD[OF i-set-mult-subset-i-set[of k div d]]*
 ⟨proof⟩

corollary *i-set-mult-div-self-i-set*:
 $I \in i\text{-set-mult } k \implies I \oslash k \in i\text{-set}$
 ⟨proof⟩

lemma *i-set-mod-0-mult-in-i-set-mult*:
 $\llbracket I \in i\text{-set}; m \bmod k = 0 \rrbracket \implies I \otimes m \in i\text{-set-mult } k$
 ⟨proof⟩

lemma *i-set-self-mult-in-i-set-mult*:
 $I \in i\text{-set} \implies I \otimes k \in i\text{-set-mult } k$
 ⟨proof⟩

lemma *i-set-mult-mod-gr0-div-not-in-i-set*:
 $\llbracket 0 < k; 0 < d; 0 < k \bmod d \rrbracket \implies \exists I \in i\text{-set-mult } k. I \oslash d \notin i\text{-set}$
 ⟨proof⟩

thm *iMOD-div-mod-gr0-not-ex[of d Suc d * k 0]*

⟨proof⟩

thm *i-set-infinite-as-iMOD*

⟨proof⟩

lemma *i-set0-mult-mod-0-div*:

$\llbracket I \in i\text{-set0-mult } k; k \bmod d = 0 \rrbracket \implies I \odot d \in i\text{-set0-mult } (k \text{ div } d)$

⟨proof⟩

corollary *i-set0-mult-mod-0-div-i-set0*:

$\llbracket I \in i\text{-set0-mult } k; k \bmod d = 0 \rrbracket \implies I \odot d \in i\text{-set0}$

⟨proof⟩

corollary *i-set0-mult-div-self-i-set0*:

$I \in i\text{-set0-mult } k \implies I \odot k \in i\text{-set0}$

⟨proof⟩

lemma *i-set0-mod-0-mult-in-i-set0-mult*:

$\llbracket I \in i\text{-set0}; m \bmod k = 0 \rrbracket \implies I \otimes m \in i\text{-set0-mult } k$

⟨proof⟩

lemma *i-set0-self-mult-in-i-set0-mult*:

$I \in i\text{-set0} \implies I \otimes k \in i\text{-set0-mult } k$

⟨proof⟩

lemma *i-set0-mult-mod-gr0-div-not-in-i-set0*:

$\llbracket 0 < k; 0 < d; 0 < k \bmod d \rrbracket \implies \exists I \in i\text{-set0-mult } k. I \odot d \notin i\text{-set0}$

⟨proof⟩

end

3 IL-TemporalOperators: Temporal logic operators on natural intervals

theory *IL-TemporalOperators*

imports *IL-IntervalOperators*

begin

Bool : some additional properties

instantiation *bool* :: {ord, zero, one, plus, times, order}

begin

definition *Zero-bool-def* [simp]: $0 \equiv \text{False}$

definition *One-bool-def* [simp]: $1 \equiv \text{True}$

definition *add-bool-def*: $a + b \equiv a \vee b$

definition *mult-bool-def*: $a * b \equiv a \wedge b$

instance ⟨proof⟩

end

value *False* < *True*

```

value True < True
value True ≤ True
thm le-bool-def

```

```

lemmas bool-op-rel-defs =
  add-bool-def
  mult-bool-def
  less-bool-def
  le-bool-def
thm bool-op-rel-defs

```

```

instance bool :: linorder
<proof>

```

```

instance bool :: wellorder
<proof>
thm wf-def
<proof>

```

```

instance bool :: comm-semiring-1
<proof>

```

3.1 Basic definitions

```

typ iT

```

```

thm Ex-def
thm All-def

```

```

term All
term Ball
lemma UNIV-nat:  $\mathbf{N} = (\text{UNIV}::\text{nat set})$ 
<proof>

```

Universal temporal operator: Always/Globally

```

definition
  iAll :: iT ⇒ (Time ⇒ bool) ⇒ bool    — Always
where
  iAll I P ≡ ∀ t∈I. P t

```

Existential temporal operator: Eventually/Finally

```

definition
  iEx :: iT ⇒ (Time ⇒ bool) ⇒ bool    — Eventually
where
  iEx-def : iEx I P ≡ ∃ t∈I. P t

```

syntax (*xsymbols*)

-*iAll* :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \square - \cdot / -) [0, 0, 10] 10)$
 -*iEx* :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \diamond - \cdot / -) [0, 0, 10] 10)$

syntax (*HTML output*)

-*iAll* :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \square - \cdot / -) [0, 0, 10] 10)$
 -*iEx* :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \diamond - \cdot / -) [0, 0, 10] 10)$

translations

$\square t I. P \equiv CONST\ iAll\ I\ (\lambda t. P)$
 $\diamond t I. P \equiv CONST\ iEx\ I\ (\lambda t. P)$

Future temporal operator: Next

definition

iNext :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Next

where

iNext *t0 I P* $\equiv P\ (inext\ t0\ I)$

Past temporal operator: Last/Previous

definition

iLast :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Last

where

iLast *t0 I P* $\equiv P\ (iprev\ t0\ I)$

syntax (*xsymbols*)

-*iNext* :: $Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \circ - \cdot / -) [0, 0, 10] 10)$
 -*iLast* :: $Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \ominus - \cdot / -) [0, 0, 10] 10)$

syntax (*HTML output*)

-*iNext* :: $Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \circ - \cdot / -) [0, 0, 10] 10)$
 -*iLast* :: $Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \ominus - \cdot / -) [0, 0, 10] 10)$

translations

$\circ t t0 I. P \equiv CONST\ iNext\ t0\ I\ (\lambda t. P)$
 $\ominus t t0 I. P \equiv CONST\ iLast\ t0\ I\ (\lambda t. P)$

lemma $\circ t 10 [0..]. (t + 10 > 10)$

<proof>

The following versions of Next and Last operator differ in the cases where no next/previous element exists or specified time point is not in interval: the weak versions return *True* and the strong versions return *False*.

definition

iNextWeak :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Weak Next

where

iNextWeak *t0 I P* $\equiv (\square t \{inext\ t0\ I\} \downarrow > t0. P\ t)$

definition

iNextStrong :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Strong Next

where

$$iNextStrong\ t0\ I\ P \equiv (\diamond t \{inext\ t0\ I\} \downarrow > t0. P\ t)$$

definition

$$iLastWeak \quad ::\ Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \quad \text{— Weak Last}$$

where

$$iLastWeak\ t0\ I\ P \equiv (\square t \{iprev\ t0\ I\} \downarrow < t0. P\ t)$$

definition

$$iLastStrong \quad ::\ Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \quad \text{— Strong Last}$$

where

$$iLastStrong\ t0\ I\ P \equiv (\diamond t \{iprev\ t0\ I\} \downarrow < t0. P\ t)$$

syntax (*xsymbols*)

$$\begin{aligned} -iNextWeak \quad &::\ Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \bigcirc_W \text{---}/ -) \\ &[0, 0, 10] 10) \end{aligned}$$

$$\begin{aligned} -iNextStrong \quad &::\ Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \bigcirc_S \text{---}/ -) \\ &[0, 0, 10] 10) \end{aligned}$$

$$\begin{aligned} -iLastWeak \quad &::\ Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \ominus_W \text{---}/ -) \\ &[0, 0, 10] 10) \end{aligned}$$

$$\begin{aligned} -iLastStrong \quad &::\ Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \ominus_S \text{---}/ -) \\ &[0, 0, 10] 10) \end{aligned}$$

syntax (*HTML output*)

$$\begin{aligned} -iNextWeak \quad &::\ Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \bigcirc_W \text{---}/ -) \\ &[0, 0, 10] 10) \end{aligned}$$

$$\begin{aligned} -iNextStrong \quad &::\ Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \bigcirc_S \text{---}/ -) \\ &[0, 0, 10] 10) \end{aligned}$$

$$\begin{aligned} -iLastWeak \quad &::\ Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \ominus_W \text{---}/ -) \\ &[0, 0, 10] 10) \end{aligned}$$

$$\begin{aligned} -iLastStrong \quad &::\ Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \ ((\exists \ominus_S \text{---}/ -) \\ &[0, 0, 10] 10) \end{aligned}$$

translations

$$\bigcirc_W\ t\ t0\ I. P \equiv CONST\ iNextWeak\ t0\ I\ (\lambda t. P)$$

$$\bigcirc_S\ t\ t0\ I. P \equiv CONST\ iNextStrong\ t0\ I\ (\lambda t. P)$$

$$\ominus_W\ t\ t0\ I. P \equiv CONST\ iLastWeak\ t0\ I\ (\lambda t. P)$$

$$\ominus_S\ t\ t0\ I. P \equiv CONST\ iLastStrong\ t0\ I\ (\lambda t. P)$$

Some examples for Next and Last operator

$$\begin{aligned} \text{lemma } \bigcirc\ t\ 5\ [0\dots,10]. ([0::int,10,20,30,40,50,60,70,80,90] ! t < 80) \\ \langle proof \rangle \end{aligned}$$

$$\begin{aligned} \text{lemma } \ominus\ t\ 5\ [0\dots,10]. ([0::int,10,20,30,40,50,60,70,80,90] ! t < 80) \\ \langle proof \rangle \end{aligned}$$

Temporal Until operator

definition

$$iUntil \quad ::\ iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool \quad \text{— Until}$$

where

$$iUntil\ I\ P\ Q \equiv \diamond t\ I. Q\ t \wedge (\square t' (I \downarrow < t). P\ t')$$

Temporal Since operator (past operator corresponding to Until)

definition

$$i\text{Since} \quad :: \quad iT \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow \text{bool} \quad \text{— Since}$$
where

$$i\text{Since } I P Q \equiv \diamond t I. Q t \wedge (\Box t' (I \downarrow > t). P t')$$
syntax (*xsymbols*)
$$\begin{aligned} -i\text{Until} &:: \text{Time} \Rightarrow \text{Time} \Rightarrow iT \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow \text{bool} \\ &((\cdot / - (\exists \mathcal{U} - \cdot) / -) [10, 0, 0, 0, 10] 10) \end{aligned}$$

$$\begin{aligned} -i\text{Since} &:: \text{Time} \Rightarrow \text{Time} \Rightarrow iT \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow \text{bool} \\ &((\cdot / - (\exists \mathcal{S} - \cdot) / -) [10, 0, 0, 0, 10] 10) \end{aligned}$$
translations

$$P. t \mathcal{U} t' I. Q \Rightarrow \text{CONST } i\text{Until } I (\lambda t. P) (\lambda t'. Q)$$

$$P. t \mathcal{S} t' I. Q \Rightarrow \text{CONST } i\text{Since } I (\lambda t. P) (\lambda t'. Q)$$
definition

$$i\text{WeakUntil} \quad :: \quad iT \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow \text{bool} \quad \text{— Weak Until/Wating-for/Unless}$$
where

$$i\text{WeakUntil } I P Q \equiv (\Box t I. P t) \vee (\diamond t I. Q t \wedge (\Box t' (I \downarrow < t). P t'))$$
definition

$$i\text{WeakSince} \quad :: \quad iT \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow \text{bool} \quad \text{— Weak Since/Back-to}$$
where

$$i\text{WeakSince } I P Q \equiv (\Box t I. P t) \vee (\diamond t I. Q t \wedge (\Box t' (I \downarrow > t). P t'))$$
syntax (*xsymbols*)
$$\begin{aligned} -i\text{WeakUntil} &:: \text{Time} \Rightarrow \text{Time} \Rightarrow iT \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow \text{bool} \\ &((\cdot / - (\exists \mathcal{W} - \cdot) / -) [10, 0, 0, 0, 10] 10) \end{aligned}$$

$$\begin{aligned} -i\text{WeakSince} &:: \text{Time} \Rightarrow \text{Time} \Rightarrow iT \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow \text{bool} \\ &((\cdot / - (\exists \mathcal{B} - \cdot) / -) [10, 0, 0, 0, 10] 10) \end{aligned}$$
translations

$$P. t \mathcal{W} t' I. Q \Rightarrow \text{CONST } i\text{WeakUntil } I (\lambda t. P) (\lambda t'. Q)$$

$$P. t \mathcal{B} t' I. Q \Rightarrow \text{CONST } i\text{WeakSince } I (\lambda t. P) (\lambda t'. Q)$$
definition

$$i\text{Release} \quad :: \quad iT \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow \text{bool} \quad \text{— Release}$$
where

$$i\text{Release-def} : i\text{Release } I P Q \equiv (\Box t I. Q t) \vee (\diamond t I. P t \wedge (\Box t' (I \downarrow \leq t). Q t'))$$
definition

$$i\text{Trigger} \quad :: \quad iT \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow (\text{Time} \Rightarrow \text{bool}) \Rightarrow \text{bool} \quad \text{— Trigger}$$
where

$$i\text{Trigger-def} : i\text{Trigger } I P Q \equiv (\Box t I. Q t) \vee (\diamond t I. P t \wedge (\Box t' (I \downarrow \geq t). Q t'))$$

$t')$)

syntax (*xsymbols*)

$-iRelease :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$
 $((./ - (\mathcal{R} - -). / -) [10, 0, 0, 0, 10] 10)$

$-iTrigger :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$
 $((./ - (\mathcal{T} - -). / -) [10, 0, 0, 0, 10] 10)$

translations

$P. t \mathcal{R} t' I. Q \equiv CONST\ iRelease\ I\ (\lambda t. P)\ (\lambda t'. Q)$

$P. t \mathcal{T} t' I. Q \equiv CONST\ iTrigger\ I\ (\lambda t. P)\ (\lambda t'. Q)$

lemmas *iTL-Next-defs* =

iNext-def *iLast-def*

iNextWeak-def *iLastWeak-def*

iNextStrong-def *iLastStrong-def*

lemmas *iTL-defs* =

iAll-def *iEx-def*

iUntil-def *iSince-def*

iWeakUntil-def *iWeakSince-def*

iRelease-def *iTrigger-def*

iTL-Next-defs

$\langle ML \rangle$

term $\square t I. P t$

term $\diamond t I. P t$

term $P1\ t1. t1\ \mathcal{U}\ t2\ I. P2\ t2$

term $P1\ t1. t1\ \mathcal{S}\ t2\ I. P2\ t2$

$\langle ML \rangle$

term $P\ t1. t1\ \mathcal{U}\ t2\ I. Q\ t2$

term $P\ t1. t1\ \mathcal{S}\ t2\ I. Q\ t2$

$\langle ML \rangle$

term $P\ t1. t1\ \mathcal{W}\ t2\ I. Q\ t2$

term $P\ t1. t1\ \mathcal{B}\ t2\ I. Q\ t2$

$\langle ML \rangle$

term $P t1. t1 \mathcal{R} t2 I. Q t2$

term $P t1. t1 \mathcal{T} t2 I. Q t2$

3.2 Basic lemmata for temporal operators

3.2.1 Intro/elim rules

thm *bexI rev-bexI*

lemma

iexI[intro]: $\llbracket P t; t \in I \rrbracket \Longrightarrow \diamond t I. P t$ **and**

rev-iexI[intro?]: $\llbracket t \in I; P t \rrbracket \Longrightarrow \diamond t I. P t$ **and**

iexE[elim!]: $\llbracket \diamond t I. P t; \bigwedge t. \llbracket t \in I; P t \rrbracket \Longrightarrow Q \rrbracket \Longrightarrow Q$

<proof>

lemma

iallI[intro!]: $(\bigwedge t. t \in I \Longrightarrow P t) \Longrightarrow \square t I. P t$ **and**

ispec[dest?]: $\llbracket \square t I. P t; t \in I \rrbracket \Longrightarrow P t$ **and**

iallE[elim]: $\llbracket \square t I. P t; P t \Longrightarrow Q; t \notin I \Longrightarrow Q \rrbracket \Longrightarrow Q$

<proof>

lemma

iuntilI[intro]:

$\llbracket Q t; (\bigwedge t'. t' \in I \downarrow < t \Longrightarrow P t'); t \in I \rrbracket \Longrightarrow P t'. t' \mathcal{U} t I. Q t$ **and**

rev-iuntilI[intro?]:

$\llbracket t \in I; Q t; (\bigwedge t'. t' \in I \downarrow < t \Longrightarrow P t') \rrbracket \Longrightarrow P t'. t' \mathcal{U} t I. Q t$

<proof>

lemma

iuntilE[elim]:

$\llbracket P' t'. t' \mathcal{U} t I. P t; \bigwedge t. \llbracket t \in I; P t \rrbracket \Longrightarrow Q \rrbracket \Longrightarrow Q$

<proof>

thm *iSince-def*

lemma

isinceI[intro]:

$\llbracket Q t; (\bigwedge t'. t' \in I \downarrow > t \Longrightarrow P t'); t \in I \rrbracket \Longrightarrow P t'. t' \mathcal{S} t I. Q t$ **and**

rev-isinceI[intro?]:

$\llbracket t \in I; Q t; (\bigwedge t'. t' \in I \downarrow > t \Longrightarrow P t') \rrbracket \Longrightarrow P t'. t' \mathcal{S} t I. Q t$

<proof>

lemma

isinceE[elim]:

$\llbracket P' t'. t' \mathcal{S} t I. P t; \bigwedge t. \llbracket t \in I; P t \rrbracket \Longrightarrow Q \rrbracket \Longrightarrow Q$

<proof>

3.2.2 Rewrite rules for trivial simplification

thm *ball-triv*

lemma *iall-triv[simp]*: $(\square t I. P) = ((\exists t. t \in I) \longrightarrow P)$

<proof>

thm *bex-triv*

lemma *iex-triv[simp]*: $(\diamond t I. P) = ((\exists t. t \in I) \wedge P)$

<proof>

lemma *iex-conjL1*:

$$\begin{aligned} & (\diamond t1 I1. (P1 t1 \wedge (\diamond t2 I2. P2 t1 t2))) = \\ & (\diamond t1 I1. \diamond t2 I2. P1 t1 \wedge P2 t1 t2) \end{aligned}$$

<proof>

lemma *iex-conjR1*:

$$\begin{aligned} & (\diamond t1 I1. ((\diamond t2 I2. P2 t1 t2) \wedge P1 t1)) = \\ & (\diamond t1 I1. \diamond t2 I2. P2 t1 t2 \wedge P1 t1) \end{aligned}$$

<proof>

lemma *iex-conjL2*:

$$\begin{aligned} & (\diamond t1 I1. (P1 t1 \wedge (\diamond t2 (I2 t1). P2 t1 t2))) = \\ & (\diamond t1 I1. \diamond t2 (I2 t1). P1 t1 \wedge P2 t1 t2) \end{aligned}$$

<proof>

lemma *iex-conjR2*:

$$\begin{aligned} & (\diamond t1 I1. ((\diamond t2 (I2 t1). P2 t1 t2) \wedge P1 t1)) = \\ & (\diamond t1 I1. \diamond t2 (I2 t1). P2 t1 t2 \wedge P1 t1) \end{aligned}$$

<proof>

lemma *iex-commute*:

$$\begin{aligned} & (\diamond t1 I1. \diamond t2 I2. P t1 t2) = \\ & (\diamond t2 I2. \diamond t1 I1. P t1 t2) \end{aligned}$$

<proof>

lemma *iall-conjL1*:

$$\begin{aligned} & I2 \neq \{\} \implies \\ & (\Box t1 I1. (P1 t1 \wedge (\Box t2 I2. P2 t1 t2))) = \\ & (\Box t1 I1. \Box t2 I2. P1 t1 \wedge P2 t1 t2) \end{aligned}$$

<proof>

lemma *iall-conjR1*:

$$\begin{aligned} & I2 \neq \{\} \implies \\ & (\Box t1 I1. ((\Box t2 I2. P2 t1 t2) \wedge P1 t1)) = \\ & (\Box t1 I1. \Box t2 I2. P2 t1 t2 \wedge P1 t1) \end{aligned}$$

<proof>

lemma *iall-conjL2*:

$$\begin{aligned} & \Box t1 I1. I2 t1 \neq \{\} \implies \\ & (\Box t1 I1. (P1 t1 \wedge (\Box t2 (I2 t1). P2 t1 t2))) = \\ & (\Box t1 I1. \Box t2 (I2 t1). P1 t1 \wedge P2 t1 t2) \end{aligned}$$

<proof>

lemma *iall-conjR2*:

$$\begin{aligned} & \Box t1 I1. I2 t1 \neq \{\} \implies \\ & (\Box t1 I1. ((\Box t2 (I2 t1). P2 t1 t2) \wedge P1 t1)) = \\ & (\Box t1 I1. \Box t2 (I2 t1). P2 t1 t2 \wedge P1 t1) \end{aligned}$$

<proof>

lemma *iall-commute*:

$$\begin{aligned} & (\Box t1 I1. \Box t2 I2. P t1 t2) = \\ & (\Box t2 I2. \Box t1 I1. P t1 t2) \end{aligned}$$

<proof>

lemma *iall-conj-distrib*:

$$(\Box t I. P t \wedge Q t) = ((\Box t I. P t) \wedge (\Box t I. Q t))$$

<proof>

lemma *iex-disj-distrib*:

$$(\Diamond t I. P t \vee Q t) = ((\Diamond t I. P t) \vee (\Diamond t I. Q t))$$

<proof>

lemma *iall-conj-distrib2*:

$$\begin{aligned} &(\Box t1 I1. \Box t2 (I2 t1). P t1 t2 \wedge Q t1 t2) = \\ &((\Box t1 I1. \Box t2 (I2 t1). P t1 t2) \wedge (\Box t1 I1. \Box t2 (I2 t1). Q t1 t2)) \end{aligned}$$

<proof>

lemma *iex-disj-distrib2*:

$$\begin{aligned} &(\Diamond t1 I1. \Diamond t2 (I2 t1). P t1 t2 \vee Q t1 t2) = \\ &((\Diamond t1 I1. \Diamond t2 (I2 t1). P t1 t2) \vee (\Diamond t1 I1. \Diamond t2 (I2 t1). Q t1 t2)) \end{aligned}$$

<proof>

lemma *iUntil-disj-distrib*:

$$(P t1. t1 U t2 I. (Q1 t2 \vee Q2 t2)) = ((P t1. t1 U t2 I. Q1 t2) \vee (P t1. t1 U t2 I. Q2 t2))$$

<proof>

lemma *iSince-disj-distrib*:

$$(P t1. t1 S t2 I. (Q1 t2 \vee Q2 t2)) = ((P t1. t1 S t2 I. Q1 t2) \vee (P t1. t1 S t2 I. Q2 t2))$$

<proof>

thm *iNextWeak-def*

lemma

$$iNext\text{-iff} \quad : (\bigcirc t t0 I. P t) = (\Box t [\dots 0] \oplus (inext t0 I). P t) \text{ and}$$

$$iLast\text{-iff} \quad : (\ominus t t0 I. P t) = (\Box t [\dots 0] \oplus (iprev t0 I). P t)$$

<proof>

lemma

$$iNext\text{-iEx-iff} \quad : (\bigcirc t t0 I. P t) = (\Diamond t [\dots 0] \oplus (inext t0 I). P t) \text{ and}$$

$$iLast\text{-iEx-iff} \quad : (\ominus t t0 I. P t) = (\Diamond t [\dots 0] \oplus (iprev t0 I). P t)$$

<proof>

lemma *inext-singleton-cut-greater-not-empty-iff*:

$$(\{inext t0 I\} \downarrow > t0 \neq \{\}) = (I \downarrow > t0 \neq \{\}) \wedge t0 \in I$$

<proof>

thm *inext-all-le-fix*

<proof>

lemma *iprev-singleton-cut-less-not-empty-iff*:

$$(\{iprev t0 I\} \downarrow < t0 \neq \{\}) = (I \downarrow < t0 \neq \{\}) \wedge t0 \in I$$

<proof>

thm *iprev-all-ge-fix*

<proof>

lemma *inext-singleton-cut-greater-empty-iff*:

$$(\{inext t0 I\} \downarrow > t0 = \{\}) = (I \downarrow > t0 = \{\}) \vee t0 \notin I$$

<proof>

thm *inext-singleton-cut-greater-not-empty-iff*

<proof>

lemma *iprev-singleton-cut-less-empty-iff*:

$$\{\text{iprev } t0 \ I\} \downarrow < t0 = \{\}\} = (I \downarrow < t0 = \{\} \vee t0 \notin I)$$

<proof>

thm *iprev-singleton-cut-less-not-empty-iff*

<proof>

lemma *iNextWeak-iff* : $(\bigcirc_W t \ t0 \ I. \ P \ t) = ((\bigcirc t \ t0 \ I. \ P \ t) \vee (I \downarrow > t0 = \{\}) \vee t0 \notin I)$

<proof>

thm *inext-singleton-cut-greater-empty-iff* [symmetric]

<proof>

lemma *iNextStrong-iff* : $(\bigcirc_S t \ t0 \ I. \ P \ t) = ((\bigcirc t \ t0 \ I. \ P \ t) \wedge (I \downarrow > t0 \neq \{\}) \wedge t0 \in I)$

<proof>

thm *inext-singleton-cut-greater-not-empty-iff* [symmetric]

<proof>

lemma *iLastWeak-iff* : $(\ominus_W t \ t0 \ I. \ P \ t) = ((\ominus t \ t0 \ I. \ P \ t) \vee (I \downarrow < t0 = \{\}) \vee t0 \notin I)$

<proof>

thm *iprev-singleton-cut-less-empty-iff* [symmetric]

<proof>

lemma *iLastStrong-iff* : $(\ominus_S t \ t0 \ I. \ P \ t) = ((\ominus t \ t0 \ I. \ P \ t) \wedge (I \downarrow < t0 \neq \{\}) \wedge t0 \in I)$

<proof>

thm *iprev-singleton-cut-less-not-empty-iff* [symmetric]

<proof>

lemmas *iTL-Next-iff* =

iNext-iff *iLast-iff*

iNextWeak-iff *iNextStrong-iff*

iLastWeak-iff *iLastStrong-iff*

lemma

iNext-iff-singleton : $(\bigcirc t \ t0 \ I. \ P \ t) = (\Box t \ \{\text{inext } t0 \ I\}. \ P \ t)$ **and**

iLast-iff-singleton : $(\ominus t \ t0 \ I. \ P \ t) = (\Box t \ \{\text{iprev } t0 \ I\}. \ P \ t)$

<proof>

lemmas *iNextLast-iff-singleton* = *iNext-iff-singleton* *iLast-iff-singleton*

lemma

iNext-iEx-iff-singleton : $(\bigcirc t \ t0 \ I. \ P \ t) = (\Diamond t \ \{\text{inext } t0 \ I\}. \ P \ t)$ **and**

iLast-iEx-iff-singleton : $(\ominus t \ t0 \ I. \ P \ t) = (\Diamond t \ \{\text{iprev } t0 \ I\}. \ P \ t)$

<proof>

lemma

iNextWeak-iAll-conv: $(\bigcirc_W t t0 I. P t) = (\Box t (\{inext t0 I\} \downarrow > t0). P t)$ **and**
iNextStrong-iEx-conv: $(\bigcirc_S t t0 I. P t) = (\Diamond t (\{inext t0 I\} \downarrow > t0). P t)$ **and**
iLastWeak-iAll-conv: $(\ominus_W t t0 I. P t) = (\Box t (\{iprev t0 I\} \downarrow < t0). P t)$ **and**
iLastStrong-iEx-conv: $(\ominus_S t t0 I. P t) = (\Diamond t (\{iprev t0 I\} \downarrow < t0). P t)$
 <proof>

lemma

iAll-True[simp]: $\Box t I. True$ **and**
iAll-False[simp]: $(\Box t I. False) = (I = \{\})$ **and**
iEx-True[simp]: $(\Diamond t I. True) = (I \neq \{\})$ **and**
iEx-False[simp]: $\neg (\Diamond t I. False)$
 <proof>

lemma *empty-iff-iAll-False*: $(I = \{\}) = (\Box t I. False)$ <proof>

lemma *not-empty-iff-iEx-True*: $(I \neq \{\}) = (\Diamond t I. True)$ <proof>

lemma

iNext-True: $\bigcirc t t0 I. True$ **and**
iNextWeak-True: $(\bigcirc_W t t0 I. True)$ **and**
iNext-False: $\neg (\bigcirc t t0 I. False)$ **and**
iNextStrong-False: $\neg (\bigcirc_S t t0 I. False)$
 <proof>

lemma

iNextStrong-True: $(\bigcirc_S t t0 I. True) = (I \downarrow > t0 \neq \{\} \wedge t0 \in I)$ **and**
iNextWeak-False: $(\neg (\bigcirc_W t t0 I. False)) = (I \downarrow > t0 \neq \{\} \wedge t0 \in I)$
 <proof>

lemma

iLast-True: $\ominus t t0 I. True$ **and**
iLastWeak-True: $(\ominus_W t t0 I. True)$ **and**
iLast-False: $\neg (\ominus t t0 I. False)$ **and**
iLastStrong-False: $\neg (\ominus_S t t0 I. False)$
 <proof>

lemma

iLastStrong-True: $(\ominus_S t t0 I. True) = (I \downarrow < t0 \neq \{\} \wedge t0 \in I)$ **and**
iLastWeak-False: $(\neg (\ominus_W t t0 I. False)) = (I \downarrow < t0 \neq \{\} \wedge t0 \in I)$
 <proof>

lemma *iUntil-True-left[simp]*: $(True. t' \mathcal{U} t I. Q t) = (\Diamond t I. Q t)$

<proof>

lemma *iUntil-True[simp]*: $(P t'. t' \mathcal{U} t I. True) = (I \neq \{\})$

<proof>

lemma *iUntil-False-left[simp]*: $(False. t' \mathcal{U} t I. Q t) = (I \neq \{\} \wedge Q (iMin I))$

<proof>

lemma *iUntil-False[simp]*: $\neg (P t'. t' \mathcal{U} t I. False)$

$\langle proof \rangle$

lemma *iSince-True-left[simp]*: $(True. t' \mathcal{S} t I. Q t) = (\diamond t I. Q t)$

$\langle proof \rangle$

lemma *iSince-True-if*:

$(P t'. t' \mathcal{S} t I. True) =$

$(if\ finite\ I\ then\ I \neq \{\} \ else\ \diamond t1\ I. \square t2\ (I \downarrow > t1). P\ t2)$

$\langle proof \rangle$

thm *cut-greater-Max-empty*

$\langle proof \rangle$

corollary *iSince-True-finite[simp]*: $finite\ I \implies (P t'. t' \mathcal{S} t I. True) = (I \neq \{\})$

$\langle proof \rangle$

lemma *iSince-False-left[simp]*: $(False. t' \mathcal{S} t I. Q t) = (finite\ I \wedge I \neq \{\} \wedge Q (Max\ I))$

$\langle proof \rangle$

lemma *iSince-False[simp]*: $\neg (P t'. t' \mathcal{S} t I. False)$

$\langle proof \rangle$

lemma *iWeakUntil-True-left[simp]*: $True. t' \mathcal{W} t I. Q t$

$\langle proof \rangle$

lemma *iWeakUntil-True[simp]*: $P t'. t' \mathcal{W} t I. True$

$\langle proof \rangle$

lemma *iWeakUntil-False-left[simp]*: $(False. t' \mathcal{W} t I. Q t) = (I = \{\} \vee Q (iMin\ I))$

$\langle proof \rangle$

lemma *iWeakUntil-False[simp]*: $(P t'. t' \mathcal{W} t I. False) = (\square t I. P t)$

$\langle proof \rangle$

lemma *iWeakSince-True-left[simp]*: $True. t' \mathcal{B} t I. Q t$

$\langle proof \rangle$

lemma *iWeakSince-True-disj*:

$(P t'. t' \mathcal{B} t I. True) =$

$(I = \{\} \vee (\diamond t1\ I. \square t2\ (I \downarrow > t1). P\ t2))$

$\langle proof \rangle$

lemma *iWeakSince-True-finite[simp]*: $finite\ I \implies (P t'. t' \mathcal{B} t I. True)$

$\langle proof \rangle$

lemma *iWeakSince-False-left[simp]*: $(False. t' \mathcal{B} t I. Q t) = (I = \{\} \vee (finite\ I \wedge Q (Max\ I)))$

$\langle proof \rangle$

lemma *iWeakSince-False[simp]*: $(P t'. t' \mathcal{B} t I. False) = (\square t I. P t)$

$\langle proof \rangle$

lemma *iRelease-True-left[simp]*: $(True. t' \mathcal{R} t I. Q t) = (I = \{\} \vee Q (iMin\ I))$

$\langle proof \rangle$

lemma *iRelease-True[simp]*: $P t'. t' \mathcal{R} t I. True$

$\langle proof \rangle$

lemma *iRelease-False-left[simp]*: $(False. t' \mathcal{R} t I. Q t) = (\square t I. Q t)$

$\langle proof \rangle$

lemma *iRelease-False[simp]*: $(P t'. t' \mathcal{R} t I. False) = (I = \{\})$

$\langle proof \rangle$

lemma $iTrigger-True-left[simp]$: $(True. t' \mathcal{T} t I. Q t) = (I = \{\}) \vee (\diamond t1 I. \square t2 (I \downarrow \geq t1). Q t2)$

$\langle proof \rangle$

lemma $iTrigger-True[simp]$: $P t'. t' \mathcal{T} t I. True$

$\langle proof \rangle$

lemma $iTrigger-False-left[simp]$: $(False. t' \mathcal{T} t I. Q t) = (\square t I. Q t)$

$\langle proof \rangle$

lemma $iTrigger-False[simp]$: $(P t'. t' \mathcal{T} t I. False) = (I = \{\})$

$\langle proof \rangle$

lemma

$iUntil-TrueTrue[simp]$: $(True. t' \mathcal{U} t I. True) = (I \neq \{\})$ **and**

$iSince-TrueTrue[simp]$: $(True. t' \mathcal{S} t I. True) = (I \neq \{\})$ **and**

$iWeakUntil-TrueTrue[simp]$: $True. t' \mathcal{W} t I. True$ **and**

$iWeakSince-TrueTrue[simp]$: $True. t' \mathcal{B} t I. True$ **and**

$iRelease-TrueTrue[simp]$: $True. t' \mathcal{R} t I. True$ **and**

$iTrigger-TrueTrue[simp]$: $True. t' \mathcal{T} t I. True$

$\langle proof \rangle$

3.2.3 Empty sets and singletons

lemma $iAll-empty[simp]$: $\square t \{\}$. $P t$ $\langle proof \rangle$

lemma $iEx-empty[simp]$: $\neg (\diamond t \{\}). P t$ $\langle proof \rangle$

thm $iAll-empty iEx-empty$

lemma $iUntil-empty[simp]$: $\neg (P t0. t0 \mathcal{U} t1 \{\}. Q t1)$ $\langle proof \rangle$

lemma $iSince-empty[simp]$: $\neg (P t0. t0 \mathcal{S} t1 \{\}. Q t1)$ $\langle proof \rangle$

lemma $iWeakUntil-empty[simp]$: $P t0. t0 \mathcal{W} t1 \{\}. Q t1$ $\langle proof \rangle$

lemma $iWeakSince-empty[simp]$: $P t0. t0 \mathcal{B} t1 \{\}. Q t1$ $\langle proof \rangle$

lemma $iRelease-empty[simp]$: $P t0. t0 \mathcal{R} t1 \{\}. Q t1$ $\langle proof \rangle$

lemma $iTrigger-empty[simp]$: $P t0. t0 \mathcal{T} t1 \{\}. Q t1$ $\langle proof \rangle$

lemmas $iTL-empty =$

$iAll-empty iEx-empty$

$iUntil-empty iSince-empty$

$iWeakUntil-empty iWeakSince-empty$

$iRelease-empty iTrigger-empty$

lemma $iAll-singleton[simp]$: $(\square t' \{t\}. P t') = P t$ $\langle proof \rangle$

lemma $iEx-singleton[simp]$: $(\diamond t' \{t\}. P t') = P t$ $\langle proof \rangle$

lemma $iUntil-singleton[simp]$: $(P t0. t0 \mathcal{U} t1 \{t\}. Q t1) = Q t$
 $\langle proof \rangle$

lemma $iSince-singleton[simp]$: $(P t0. t0 \mathcal{S} t1 \{t\}. Q t1) = Q t$
 $\langle proof \rangle$

lemma *iWeakUntil-singleton[simp]*: $(P\ t0.\ t0\ \mathcal{W}\ t1\ \{t\}.\ Q\ t1) = (P\ t \vee Q\ t)$
 $\langle proof \rangle$

lemma *iWeakSince-singleton[simp]*: $(P\ t0.\ t0\ \mathcal{B}\ t1\ \{t\}.\ Q\ t1) = (P\ t \vee Q\ t)$
 $\langle proof \rangle$

lemma *iRelease-singleton[simp]*: $(P\ t0.\ t0\ \mathcal{R}\ t1\ \{t\}.\ Q\ t1) = Q\ t$
 $\langle proof \rangle$

lemma *iTrigger-singleton[simp]*: $(P\ t0.\ t0\ \mathcal{T}\ t1\ \{t\}.\ Q\ t1) = Q\ t$
 $\langle proof \rangle$

lemmas *iTL-singleton* =
iAll-singleton *iEx-singleton*
iUntil-singleton *iSince-singleton*
iWeakUntil-singleton *iWeakSince-singleton*
iRelease-singleton *iTrigger-singleton*

thm *iTL-singleton*

3.2.4 Conversions between temporal operators

lemma *iAll-iEx-conv*: $(\Box\ t\ I.\ P\ t) = (\neg\ (\Diamond\ t\ I.\ \neg\ P\ t))$ $\langle proof \rangle$

lemma *iEx-iAll-conv*: $(\Diamond\ t\ I.\ P\ t) = (\neg\ (\Box\ t\ I.\ \neg\ P\ t))$ $\langle proof \rangle$

lemma *not-iAll[simp]*: $(\neg\ (\Box\ t\ I.\ P\ t)) = (\Diamond\ t\ I.\ \neg\ P\ t)$ $\langle proof \rangle$

lemma *not-iEx[simp]*: $(\neg\ (\Diamond\ t\ I.\ P\ t)) = (\Box\ t\ I.\ \neg\ P\ t)$ $\langle proof \rangle$

lemma *iUntil-iEx-conv*: $(True.\ t'\ \mathcal{U}\ t\ I.\ P\ t) = (\Diamond\ t\ I.\ P\ t)$ $\langle proof \rangle$

lemma *iSince-iEx-conv*: $(True.\ t'\ \mathcal{S}\ t\ I.\ P\ t) = (\Diamond\ t\ I.\ P\ t)$ $\langle proof \rangle$

lemma *iRelease-iAll-conv*: $(False.\ t'\ \mathcal{R}\ t\ I.\ P\ t) = (\Box\ t\ I.\ P\ t)$
 $\langle proof \rangle$

lemma *iTrigger-iAll-conv*: $(False.\ t'\ \mathcal{T}\ t\ I.\ P\ t) = (\Box\ t\ I.\ P\ t)$
 $\langle proof \rangle$

lemma *iWeakUntil-iUntil-conv*: $(P\ t'.\ t'\ \mathcal{W}\ t\ I.\ Q\ t) = ((P\ t'.\ t'\ \mathcal{U}\ t\ I.\ Q\ t) \vee (\Box\ t\ I.\ P\ t))$
 $\langle proof \rangle$

lemma *iWeakSince-iSince-conv*: $(P\ t'.\ t'\ \mathcal{B}\ t\ I.\ Q\ t) = ((P\ t'.\ t'\ \mathcal{S}\ t\ I.\ Q\ t) \vee (\Box\ t\ I.\ P\ t))$
 $\langle proof \rangle$

lemma *iUntil-iWeakUntil-conv*: $(P\ t'.\ t'\ \mathcal{U}\ t\ I.\ Q\ t) = ((P\ t'.\ t'\ \mathcal{W}\ t\ I.\ Q\ t) \wedge (\Diamond\ t\ I.\ Q\ t))$
 $\langle proof \rangle$

lemma *iSince-iWeakSince-conv*: $(P\ t'.\ t'\ \mathcal{S}\ t\ I.\ Q\ t) = ((P\ t'.\ t'\ \mathcal{B}\ t\ I.\ Q\ t) \wedge (\Diamond\ t\ I.\ Q\ t))$
 $\langle proof \rangle$

thm *iRelease-def* *iWeakUntil-def*

lemma *iRelease-iWeakUntil-conv*: $(P t'. t' \mathcal{R} t I. Q t) = (Q t'. t' \mathcal{W} t I. (Q t \wedge P t))$

<proof>

lemma *iRelease-iUntil-conv*: $(P t'. t' \mathcal{R} t I. Q t) = ((\Box t I. Q t) \vee (Q t'. t' \mathcal{U} t I. (Q t \wedge P t)))$

<proof>

lemma *iTrigger-iWeakSince-conv*: $(P t'. t' \mathcal{T} t I. Q t) = (Q t'. t' \mathcal{B} t I. (Q t \wedge P t))$

<proof>

lemma *iTrigger-iSince-conv*: $(P t'. t' \mathcal{T} t I. Q t) = ((\Box t I. Q t) \vee (Q t'. t' \mathcal{S} t I. (Q t \wedge P t)))$

<proof>

lemma *iRelease-not-iUntil-conv*: $(P t'. t' \mathcal{R} t I. Q t) = (\neg (\neg P t'. t' \mathcal{U} t I. \neg Q t))$

<proof>

lemma *iUntil-not-iRelease-conv*: $(P t'. t' \mathcal{U} t I. Q t) = (\neg (\neg P t'. t' \mathcal{R} t I. \neg Q t))$

<proof>

The Trigger operator \mathcal{T} is a past operator, so that it is used for time intervals, that are bounded by a current time point, and thus are finite. For an infinite interval the stated relation to the Since operator \mathcal{S} would not be fulfilled.

lemma *iTrigger-not-iSince-conv*: *finite I* $\implies (P t'. t' \mathcal{T} t I. Q t) = (\neg (\neg P t'. t' \mathcal{S} t I. \neg Q t))$

<proof>

thm *not-greater-Max*[rotated 1]

<proof>

lemma *iSince-not-iTrigger-conv*: *finite I* $\implies (P t'. t' \mathcal{S} t I. Q t) = (\neg (\neg P t'. t' \mathcal{T} t I. \neg Q t))$

<proof>

lemma *not-iUntil*:

$$\begin{aligned} & (\neg (P t1. t1 \mathcal{U} t2 I. Q t2)) = \\ & (\Box t I. (Q t \longrightarrow (\Diamond t' (I \downarrow < t). \neg P t'))) \end{aligned}$$

<proof>

lemma *not-iSince*:

$$\begin{aligned} & (\neg (P t1. t1 \mathcal{S} t2 I. Q t2)) = \\ & (\Box t I. (Q t \longrightarrow (\Diamond t' (I \downarrow > t). \neg P t'))) \end{aligned}$$

<proof>

lemma *iWeakUntil-conj-iUntil-conv*:

$$(P t1. t1 \mathcal{W} t2 I. (P t2 \wedge Q t2)) = (\neg (\neg Q t1. t1 \mathcal{U} t2 I. \neg P t2))$$

<proof>

lemma *iUntil-disj-iUntil-conv*:

$$(P t1 \vee Q t1. t1 \mathcal{U} t2 I. Q t2) =$$

$(P t1. t1 \mathcal{U} t2 I. Q t2)$
 ⟨proof⟩
thm *subsetD[OF - iMinI]*
 ⟨proof⟩
lemma *iWeakUntil-disj-iWeakUntil-conv:*
 $(P t1 \vee Q t1. t1 \mathcal{W} t2 I. Q t2) =$
 $(P t1. t1 \mathcal{W} t2 I. Q t2)$
 ⟨proof⟩
lemma *iWeakUntil-iUntil-conj-conv:*
 $(P t1. t1 \mathcal{W} t2 I. Q t2) =$
 $(\neg (\neg Q t1. t1 \mathcal{U} t2 I. (\neg P t2 \wedge \neg Q t2)))$
 ⟨proof⟩

Negation and temporal operators

thm *iNext-iff*

lemma

not-iNext[simp]: $(\neg (\bigcirc t t0 I. P t)) = (\bigcirc t t0 I. \neg P t)$ **and**
not-iNextWeak[simp]: $(\neg (\bigcirc_W t t0 I. P t)) = (\bigcirc_S t t0 I. \neg P t)$ **and**
not-iNextStrong[simp]: $(\neg (\bigcirc_S t t0 I. P t)) = (\bigcirc_W t t0 I. \neg P t)$ **and**
not-iLast[simp]: $(\neg (\ominus t t0 I. P t)) = (\ominus t t0 I. \neg P t)$ **and**
not-iLastWeak[simp]: $(\neg (\ominus_W t t0 I. P t)) = (\ominus_S t t0 I. \neg P t)$ **and**
not-iLastStrong[simp]: $(\neg (\ominus_S t t0 I. P t)) = (\ominus_W t t0 I. \neg P t)$
 ⟨proof⟩

thm *not-iUntil*

thm *not-iSince*

lemma *not-iWeakUntil:*

$(\neg (P t1. t1 \mathcal{W} t2 I. Q t2)) =$
 $((\Box t I. (Q t \longrightarrow (\Diamond t' (I \downarrow < t). \neg P t'))) \wedge (\Diamond t I. \neg P t))$
 ⟨proof⟩

lemma *not-iWeakSince:*

$(\neg (P t1. t1 \mathcal{B} t2 I. Q t2)) =$
 $((\Box t I. (Q t \longrightarrow (\Diamond t' (I \downarrow > t). \neg P t'))) \wedge (\Diamond t I. \neg P t))$
 ⟨proof⟩

lemma *not-iRelease:*

$(\neg (P t'. t' \mathcal{R} t I. Q t)) =$
 $((\Diamond t I. \neg Q t) \wedge (\Box t I. P t \longrightarrow (\Diamond t I \downarrow \leq t. \neg Q t)))$
 ⟨proof⟩

lemma *not-iRelease-iUntil-conv:*

$(\neg (P t'. t' \mathcal{R} t I. Q t)) = (\neg P t'. t' \mathcal{U} t I. \neg Q t)$
 ⟨proof⟩

lemma *not-iTrigger:*

$(\neg (P t'. t' \mathcal{T} t I. Q t)) =$
 $((\Diamond t I. \neg Q t) \wedge (\Box t I. \neg P t \vee (\Diamond t I \downarrow \geq t. \neg Q t)))$
 ⟨proof⟩

lemma *not-iTrigger-iSince-conv:*

$finite I \implies (\neg (P t'. t' \mathcal{T} t I. Q t)) = (\neg P t'. t' \mathcal{S} t I. \neg Q t)$

<proof>

3.2.5 Some implication results

lemma *all-imp-iall*: $\forall x. P x \implies \Box t I. P t$ *<proof>*

lemma *bex-imp-lex*: $\Diamond t I. P t \implies \exists x. P x$ *<proof>*

lemma *iAll-imp-iEx*: $I \neq \{\}$ $\implies \Box t I. P t \implies \Diamond t I. P t$ *<proof>*

lemma *i-set-iAll-imp-iEx*: $I \in i\text{-set} \implies \Box t I. P t \implies \Diamond t I. P t$
<proof>

lemmas *iT-iAll-imp-iEx = iT-not-empty[THEN iAll-imp-iEx]*

thm *iT-iAll-imp-iEx*

lemma *iUntil-imp-iEx*: $P t1. t1 \mathcal{U} t2 I. Q t2 \implies \Diamond t I. Q t$

<proof>

lemma *iSince-imp-iEx*: $P t1. t1 \mathcal{S} t2 I. Q t2 \implies \Diamond t I. Q t$

<proof>

thm *ball-subset-imp-ball*

lemma *iall-subset-imp-iall*: $\llbracket \Box t B. P t; A \subseteq B \rrbracket \implies \Box t A. P t$

<proof>

thm *bex-subset-imp-bex*

lemma *iex-subset-imp-iex*: $\llbracket \Diamond t A. P t; A \subseteq B \rrbracket \implies \Diamond t B. P t$

<proof>

lemma *iall-mp*: $\llbracket \Box t I. P t \longrightarrow Q t; \Box t I. P t \rrbracket \implies \Box t I. Q t$ *<proof>*

lemma *iex-mp*: $\llbracket \Box t I. P t \longrightarrow Q t; \Diamond t I. P t \rrbracket \implies \Diamond t I. Q t$ *<proof>*

lemma *iUntil-imp*:

$\llbracket P1 t1. t1 \mathcal{U} t2 I. Q t2; \Box t I. P1 t \longrightarrow P2 t \rrbracket \implies P2 t1. t1 \mathcal{U} t2 I. Q t2$

<proof>

lemma *iSince-imp*:

$\llbracket P1 t1. t1 \mathcal{S} t2 I. Q t2; \Box t I. P1 t \longrightarrow P2 t \rrbracket \implies P2 t1. t1 \mathcal{S} t2 I. Q t2$

<proof>

lemma *iWeakUntil-imp*:

$\llbracket P1 t1. t1 \mathcal{W} t2 I. Q t2; \Box t I. P1 t \longrightarrow P2 t \rrbracket \implies P2 t1. t1 \mathcal{W} t2 I. Q t2$

<proof>

lemma *iWeakSince-imp*:

$\llbracket P1 t1. t1 \mathcal{B} t2 I. Q t2; \Box t I. P1 t \longrightarrow P2 t \rrbracket \implies P2 t1. t1 \mathcal{B} t2 I. Q t2$

<proof>

lemma *iRelease-imp*:

$\llbracket P1 t1. t1 \mathcal{R} t2 I. Q t2; \Box t I. P1 t \longrightarrow P2 t \rrbracket \implies P2 t1. t1 \mathcal{R} t2 I. Q t2$

<proof>

lemma *iTrigger-imp*:

$\llbracket P1 t1. t1 \mathcal{T} t2 I. Q t2; \Box t I. P1 t \longrightarrow P2 t \rrbracket \implies P2 t1. t1 \mathcal{T} t2 I. Q t2$

<proof>

lemma *iMin-imp-iUntil*:

$$\llbracket I \neq \{\}; Q (iMin I) \rrbracket \implies P t'. t' \mathcal{U} t I. Q t$$

<proof>

lemma *Max-imp-iSince*:

$$\llbracket finite I; I \neq \{\}; Q (Max I) \rrbracket \implies P t'. t' \mathcal{S} t I. Q t$$

<proof>

3.2.6 Congruence rules for temporal operators' predicates

thm *arg-cong*

lemma *iAll-cong*: $\Box t I. f t = g t \implies (\Box t I. P (f t) t) = (\Box t I. P (g t) t)$

<proof>

lemma *iEx-cong*: $\Box t I. f t = g t \implies (\Diamond t I. P (f t) t) = (\Diamond t I. P (g t) t)$

<proof>

lemma *iUntil-cong1*:

$$\Box t I. f t = g t \implies (P (f t1) t1. t1 \mathcal{U} t2 I. Q t2) = (P (g t1) t1. t1 \mathcal{U} t2 I. Q t2)$$

<proof>

lemma *iUntil-cong2*:

$$\Box t I. f t = g t \implies (P t1. t1 \mathcal{U} t2 I. Q (f t2) t2) = (P t1. t1 \mathcal{U} t2 I. Q (g t2) t2)$$

<proof>

thm *subst*

lemma *iSince-cong1*:

$$\Box t I. f t = g t \implies (P (f t1) t1. t1 \mathcal{S} t2 I. Q t2) = (P (g t1) t1. t1 \mathcal{S} t2 I. Q t2)$$

<proof>

lemma *iSince-cong2*:

$$\Box t I. f t = g t \implies (P t1. t1 \mathcal{S} t2 I. Q (f t2) t2) = (P t1. t1 \mathcal{S} t2 I. Q (g t2) t2)$$

<proof>

thm *iAll-cong*

lemma *bex-subst*:

$$\forall x \in A. P x \longrightarrow (Q x = Q' x) \implies (\exists x \in A. P x \wedge Q x) = (\exists x \in A. P x \wedge Q' x)$$

<proof>

lemma *iEx-subst*:

$$\Box t I. P t \longrightarrow (Q t = Q' t) \implies (\Diamond t I. P t \wedge Q t) = (\Diamond t I. P t \wedge Q' t)$$

<proof>

3.2.7 Temporal operators with set unions/intersections and subsets

lemma *iAll-subset*: $\llbracket A \subseteq B; \Box t B. P t \rrbracket \implies \Box t A. P t$

<proof>

lemma *iEx-subset*: $\llbracket A \subseteq B; \Diamond t A. P t \rrbracket \implies \Diamond t B. P t$

<proof>

lemma *iUntil-append*:

$$\llbracket \text{finite } A; \text{Max } A \leq \text{iMin } B \rrbracket \Longrightarrow \\ P \ t1. \ t1 \ \mathcal{U} \ t2 \ A. \ Q \ t2 \Longrightarrow P \ t1. \ t1 \ \mathcal{U} \ t2 \ (A \cup B). \ Q \ t2$$

<proof>

thm *iEx-subset[OF Un-upper1]*

<proof>

thm *subst[OF iEx-cong, rule-format]*

<proof>

lemma *iSince-prepend*:

$$\llbracket \text{finite } A; \text{Max } A \leq \text{iMin } B \rrbracket \Longrightarrow \\ P \ t1. \ t1 \ \mathcal{S} \ t2 \ B. \ Q \ t2 \Longrightarrow P \ t1. \ t1 \ \mathcal{S} \ t2 \ (A \cup B). \ Q \ t2$$

<proof>

lemma

$$\text{iAll-union: } \llbracket \Box \ t \ A. \ P \ t; \Box \ t \ B. \ P \ t \rrbracket \Longrightarrow \Box \ t \ (A \cup B). \ P \ t \ \text{and} \\ \text{iAll-union-conv: } (\Box \ t \ A \cup B. \ P \ t) = ((\Box \ t \ A. \ P \ t) \wedge (\Box \ t \ B. \ P \ t))$$

<proof>

lemma

$$\text{iEx-union: } (\Diamond \ t \ A. \ P \ t) \vee (\Diamond \ t \ B. \ P \ t) \Longrightarrow \Diamond \ t \ (A \cup B). \ P \ t \ \text{and} \\ \text{iEx-union-conv: } (\Diamond \ t \ (A \cup B). \ P \ t) = ((\Diamond \ t \ A. \ P \ t) \vee (\Diamond \ t \ B. \ P \ t))$$

<proof>

lemma *iAll-inter: $(\Box \ t \ A. \ P \ t) \vee (\Box \ t \ B. \ P \ t) \Longrightarrow \Box \ t \ (A \cap B). \ P \ t$ *<proof>**

lemma *not-iEx-inter*:

$$\exists A \ B \ P. \ (\Diamond \ t \ A. \ P \ t) \wedge (\Diamond \ t \ B. \ P \ t) \wedge \neg (\Diamond \ t \ (A \cap B). \ P \ t)$$

<proof>

lemma

$$\text{iAll-insert: } \llbracket P \ t; \Box \ t \ I. \ P \ t \rrbracket \Longrightarrow \Box \ t' \ (\text{insert } t \ I). \ P \ t' \ \text{and} \\ \text{iAll-insert-conv: } (\Box \ t' \ (\text{insert } t \ I). \ P \ t') = (P \ t \wedge (\Box \ t' \ I. \ P \ t'))$$

<proof>

lemma

$$\text{iEx-insert: } \llbracket P \ t \vee (\Diamond \ t \ I. \ P \ t) \rrbracket \Longrightarrow \Diamond \ t' \ (\text{insert } t \ I). \ P \ t' \ \text{and} \\ \text{iEx-insert-conv: } (\Diamond \ t' \ (\text{insert } t \ I). \ P \ t') = (P \ t \vee (\Diamond \ t' \ I. \ P \ t'))$$

<proof>

3.3 Further results for temporal operators

thm *Collect-minI-ex*

thm *Collect-minI-ex-cut*

lemma *Collect-minI-iEx: $\Diamond \ t \ I. \ P \ t \Longrightarrow \Diamond \ t \ I. \ P \ t \wedge (\Box \ t' \ (I \downarrow < t). \neg P \ t')$*

<proof>

thm *iUntil-def*

lemma *iUntil-disj-conv1*:

$$I \neq \{\} \Longrightarrow \\ (P \ t'. \ t' \ \mathcal{U} \ t \ I. \ Q \ t) = (Q \ (\text{iMin } I) \vee (P \ t'. \ t' \ \mathcal{U} \ t \ I. \ Q \ t \wedge \text{iMin } I < t))$$

<proof>

lemma *iSince-disj-conv1*:

$\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies$
 $(P t'. t' \mathcal{S} t I. Q t) = (Q (\text{Max } I) \vee (P t'. t' \mathcal{S} t I. Q t \wedge t < \text{Max } I))$
 <proof>

lemma *iUntil-next*:

$I \neq \{\} \implies$
 $(P t'. t' \mathcal{U} t I. Q t) =$
 $(Q (\text{iMin } I) \vee (P (\text{iMin } I) \wedge (P t'. t' \mathcal{U} t (I \downarrow > (\text{iMin } I)). Q t)))$
 <proof>

lemma *iSince-prev*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies$

$(P t'. t' \mathcal{S} t I. Q t) =$
 $(Q (\text{Max } I) \vee (P (\text{Max } I) \wedge (P t'. t' \mathcal{S} t (I \downarrow < \text{Max } I). Q t)))$
 <proof>

lemma *iNext-induct-rule*:

$\llbracket P (\text{iMin } I); \square t I. (P t \longrightarrow (\bigcirc t' t I. P t')) \rrbracket \implies P t$
 <proof>

lemma *iNext-induct*:

$\llbracket P (\text{iMin } I); \square t I. (P t \longrightarrow (\bigcirc t' t I. P t')) \rrbracket \implies \square t I. P t$
 <proof>

lemma *iLast-induct-rule*:

$\llbracket P (\text{Max } I); \square t I. (P t \longrightarrow (\ominus t' t I. P t')) \rrbracket \implies P t$
 <proof>

lemma *iLast-induct*:

$\llbracket P (\text{Max } I); \square t I. (P t \longrightarrow (\ominus t' t I. P t')) \rrbracket \implies \square t I. P t$
 <proof>

lemma *iUntil-conj-not*: $((P t1 \wedge \neg Q t1). t1 \mathcal{U} t2 I. Q t2) = (P t1. t1 \mathcal{U} t2 I. Q t2)$

<proof>

thm *iMin-le*

<proof>

thm *subsetD[OF - iMinI]*

<proof>

lemma *iWeakUntil-conj-not*: $((P t1 \wedge \neg Q t1). t1 \mathcal{W} t2 I. Q t2) = (P t1. t1 \mathcal{W} t2 I. Q t2)$

<proof>

lemma *iSince-conj-not*: *finite I* \implies

$((P t1 \wedge \neg Q t1). t1 \mathcal{S} t2 I. Q t2) = (P t1. t1 \mathcal{S} t2 I. Q t2)$

<proof>

thm *MaxI2*

<proof>

thm *not-greater-Max*

⟨proof⟩
thm *subsetD[OF - MaxI]*
 ⟨proof⟩
lemma *iWeakSince-conj-not: finite I* \implies
 $((P\ t1 \wedge \neg Q\ t1).\ t1\ \mathcal{B}\ t2\ I.\ Q\ t2) = (P\ t1.\ t1\ \mathcal{B}\ t2\ I.\ Q\ t2)$
 ⟨proof⟩

lemma *iNextStrong-imp-iNextWeak*: $(\bigcirc_S\ t\ t0\ I.\ P\ t) \longrightarrow (\bigcirc_W\ t\ t0\ I.\ P\ t)$
 ⟨proof⟩
lemma *iLastStrong-imp-iLastWeak*: $(\ominus_S\ t\ t0\ I.\ P\ t) \longrightarrow (\ominus_W\ t\ t0\ I.\ P\ t)$
 ⟨proof⟩

lemma *infin-imp-iNextWeak-iNextStrong-eq-iNext*:
 $\llbracket\ \text{infinite } I; t0 \in I\ \rrbracket \implies$
 $((\bigcirc_W\ t\ t0\ I.\ P\ t) = (\bigcirc\ t\ t0\ I.\ P\ t)) \wedge ((\bigcirc_S\ t\ t0\ I.\ P\ t) = (\bigcirc\ t\ t0\ I.\ P\ t))$
 ⟨proof⟩

lemma *infin-imp-iNextWeak-eq-iNext*: $\llbracket\ \text{infinite } I; t0 \in I\ \rrbracket \implies (\bigcirc_W\ t\ t0\ I.\ P\ t) = (\bigcirc\ t\ t0\ I.\ P\ t)$
 ⟨proof⟩

lemma *infin-imp-iNextStrong-eq-iNext*: $\llbracket\ \text{infinite } I; t0 \in I\ \rrbracket \implies (\bigcirc_S\ t\ t0\ I.\ P\ t) = (\bigcirc\ t\ t0\ I.\ P\ t)$
 ⟨proof⟩

lemma *infin-imp-iNextStrong-eq-iNextWeak*: $\llbracket\ \text{infinite } I; t0 \in I\ \rrbracket \implies (\bigcirc_S\ t\ t0\ I.\ P\ t) = (\bigcirc_W\ t\ t0\ I.\ P\ t)$
 ⟨proof⟩

lemma
not-in-iNext-eq: $t0 \notin I \implies (\bigcirc\ t\ t0\ I.\ P\ t) = (P\ t0)$ **and**
not-in-iLast-eq: $t0 \notin I \implies (\ominus\ t\ t0\ I.\ P\ t) = (P\ t0)$
 ⟨proof⟩

lemma
not-in-iNextWeak-eq: $t0 \notin I \implies (\bigcirc_W\ t\ t0\ I.\ P\ t) = (P\ t0)$ **and**
not-in-iLastWeak-eq: $t0 \notin I \implies (\ominus_W\ t\ t0\ I.\ P\ t) = (P\ t0)$
 ⟨proof⟩

lemma
not-in-iNextStrong-eq: $t0 \notin I \implies \neg (\bigcirc_S\ t\ t0\ I.\ P\ t) = \neg (P\ t0)$ **and**
not-in-iLastStrong-eq: $t0 \notin I \implies \neg (\ominus_S\ t\ t0\ I.\ P\ t) = \neg (P\ t0)$
 ⟨proof⟩

lemma
iNext-UNIV: $(\bigcirc\ t\ t0\ UNIV.\ P\ t) = P\ (Suc\ t0)$ **and**
iNextWeak-UNIV: $(\bigcirc_W\ t\ t0\ UNIV.\ P\ t) = P\ (Suc\ t0)$ **and**
iNextStrong-UNIV: $(\bigcirc_S\ t\ t0\ UNIV.\ P\ t) = P\ (Suc\ t0)$
 ⟨proof⟩

lemma
iLast-UNIV: $(\ominus\ t\ t0\ UNIV.\ P\ t) = P\ (t0 - Suc\ 0)$ **and**

iLastWeak-UNIV: $(\ominus_W t t0 \text{ UNIV}. P t) = (\text{if } 0 < t0 \text{ then } P (t0 - \text{Suc } 0) \text{ else True})$ **and**

iLastStrong-UNIV: $(\ominus_S t t0 \text{ UNIV}. P t) = (\text{if } 0 < t0 \text{ then } P (t0 - \text{Suc } 0) \text{ else False})$
 ⟨proof⟩

lemmas *iTL-Next-UNIV* =
iNext-UNIV iNextWeak-UNIV iNextStrong-UNIV
iLast-UNIV iLastWeak-UNIV iLastStrong-UNIV

lemma *inext-nth-iNext-Suc*: $(\bigcirc t (I \rightarrow n) I. P t) = P (I \rightarrow \text{Suc } n)$
 ⟨proof⟩

lemma *iprev-nth-iLast-Suc*: $(\ominus t (I \leftarrow n) I. P t) = P (I \leftarrow \text{Suc } n)$
 ⟨proof⟩

3.4 Temporal operators and arithmetic interval operators

Shifting intervals through addition and subtraction of constants. Mirroring intervals through subtraction of intervals from constants. Expanding and compressing intervals through multiplication and division by constants.

Always operator

lemma *iT-Plus-iAll-conv*: $(\Box t I \oplus k. P t) = (\Box t I. P (t + k))$
 ⟨proof⟩

lemma *iT-Mult-iAll-conv*: $(\Box t I \otimes k. P t) = (\Box t I. P (t * k))$
 ⟨proof⟩

lemma *iT-Plus-neg-iAll-conv*: $(\Box t I \oplus - k. P t) = (\Box t (I \downarrow \geq k). P (t - k))$
 ⟨proof⟩

lemma *iT-Minus-iAll-conv*: $(\Box t k \ominus I. P t) = (\Box t (I \downarrow \leq k). P (k - t))$
 ⟨proof⟩

lemma *iT-Div-iAll-conv*: $(\Box t I \oslash k. P t) = (\Box t I. P (t \text{ div } k))$
 ⟨proof⟩

lemmas *iT-arith-iAll-conv* =
iT-Plus-iAll-conv
iT-Mult-iAll-conv
iT-Plus-neg-iAll-conv
iT-Minus-iAll-conv
iT-Div-iAll-conv

Eventually operator

lemma
iT-Plus-iEx-conv: $(\Diamond t I \oplus k. P t) = (\Diamond t I. P (t + k))$ **and**
iT-Mult-iEx-conv: $(\Diamond t I \otimes k. P t) = (\Diamond t I. P (t * k))$ **and**
iT-Plus-neg-iEx-conv: $(\Diamond t I \oplus - k. P t) = (\Diamond t (I \downarrow \geq k). P (t - k))$ **and**
iT-Minus-iEx-conv: $(\Diamond t k \ominus I. P t) = (\Diamond t (I \downarrow \leq k). P (k - t))$ **and**
iT-Div-iEx-conv: $(\Diamond t I \oslash k. P t) = (\Diamond t I. P (t \text{ div } k))$
 ⟨proof⟩

Until and Since operators

lemma *iT-Plus-iUntil-conv*: $(P\ t1.\ t1\ U\ t2\ (I\ \oplus\ k).\ Q\ t2) = (P\ (t1\ +\ k).\ t1\ U\ t2\ I.\ Q\ (t2\ +\ k))$

<proof>

lemma *iT-Mult-iUntil-conv*: $(P\ t1.\ t1\ U\ t2\ (I\ \otimes\ k).\ Q\ t2) = (P\ (t1\ *\ k).\ t1\ U\ t2\ I.\ Q\ (t2\ *\ k))$

<proof>

lemma *iT-Plus-neg-iUntil-conv*: $(P\ t1.\ t1\ U\ t2\ (I\ \oplus\ -\ k).\ Q\ t2) = (P\ (t1\ -\ k).\ t1\ U\ t2\ (I\ \downarrow\geq\ k).\ Q\ (t2\ -\ k))$

<proof>

thm *i-cut-commute-disj*[of op $\downarrow<$ op $\downarrow\geq$, simplified]

<proof>

lemma *iT-Minus-iUntil-conv*: $(P\ t1.\ t1\ U\ t2\ (k\ \ominus\ I).\ Q\ t2) = (P\ (k\ -\ t1).\ t1\ S\ t2\ (I\ \downarrow\leq\ k).\ Q\ (k\ -\ t2))$

<proof>

lemma *iT-Div-iUntil-conv*: $(P\ t1.\ t1\ U\ t2\ (I\ \odot\ k).\ Q\ t2) = (P\ (t1\ \text{div}\ k).\ t1\ U\ t2\ I.\ Q\ (t2\ \text{div}\ k))$

<proof>

thm *cut-less-cut-ge-ident*[OF order-refl]

<proof>

find-theorems - < iMin - - \notin -

<proof>

thm *subsetD*[OF - iMinI-ex2]

<proof>

Until and Since operators can be converted into each other through substraction of intervals from constants

thm *iT-Minus-iUntil-conv*

lemma *iUntil-iSince-conv*:

$\llbracket\ \text{finite}\ I;\ \text{Max}\ I\ \leq\ k\ \rrbracket\ \implies$

$(P\ t1.\ t1\ U\ t2\ I.\ Q\ t2) = (P\ (k\ -\ t1).\ t1\ S\ t2\ (k\ \ominus\ I).\ Q\ (k\ -\ t2))$

<proof>

thm *iT-Minus-iUntil-conv*

<proof>

thm *iT-Minus-iUntil-conv*

<proof>

lemma *iSince-iUntil-conv*:

$\llbracket\ \text{finite}\ I;\ \text{Max}\ I\ \leq\ k\ \rrbracket\ \implies$

$(P\ t1.\ t1\ S\ t2\ I.\ Q\ t2) = (P\ (k\ -\ t1).\ t1\ U\ t2\ (k\ \ominus\ I).\ Q\ (k\ -\ t2))$

<proof>

thm *cut-less-Max-all*

<proof>

lemma *iT-Plus-iSince-conv*: $(P\ t1.\ t1\ S\ t2\ (I\ \oplus\ k).\ Q\ t2) = (P\ (t1\ +\ k).\ t1\ S\ t2\ I.\ Q\ (t2\ +\ k))$

<proof>

lemma *iT-Mult-iSince-conv*: $0 < k \implies (P\ t1.\ t1\ S\ t2\ (I\ \otimes\ k).\ Q\ t2) = (P\ (t1$

$* k). t1 \mathcal{S} t2 I. Q (t2 * k)$

$\langle proof \rangle$

lemma *iT-Plus-neg-iSince-conv*: $(P t1. t1 \mathcal{S} t2 (I \oplus - k). Q t2) = (P (t1 - k). t1 \mathcal{S} t2 (I \downarrow \geq k). Q (t2 - k))$

$\langle proof \rangle$

lemma *iT-Minus-iSince-conv*:

$(P t1. t1 \mathcal{S} t2 (k \ominus I). Q t2) = (P (k - t1). t1 \mathcal{U} t2 (I \downarrow \leq k). Q (k - t2))$

$\langle proof \rangle$

thm *iT-Minus-cut-eq[OF order-refl]*

$\langle proof \rangle$

thm *iSince-iUntil-conv*

$\langle proof \rangle$

lemma *iT-Div-iSince-conv*:

$0 < k \implies (P t1. t1 \mathcal{S} t2 (I \oslash k). Q t2) = (P (t1 \text{ div } k). t1 \mathcal{S} t2 I. Q (t2 \text{ div } k))$

$\langle proof \rangle$

thm *le-add-diff Suc-mod-le-divisor*

$\langle proof \rangle$

Weak Until and Weak Since operators

lemma *iT-Plus-iWeakUntil-conv*: $(P t1. t1 \mathcal{W} t2 (I \oplus k). Q t2) = (P (t1 + k). t1 \mathcal{W} t2 I. Q (t2 + k))$

$\langle proof \rangle$

lemma *iT-Mult-iWeakUntil-conv*: $(P t1. t1 \mathcal{W} t2 (I \otimes k). Q t2) = (P (t1 * k). t1 \mathcal{W} t2 I. Q (t2 * k))$

$\langle proof \rangle$

lemma *iT-Plus-neg-iWeakUntil-conv*: $(P t1. t1 \mathcal{W} t2 (I \oplus - k). Q t2) = (P (t1 - k). t1 \mathcal{W} t2 (I \downarrow \geq k). Q (t2 - k))$

$\langle proof \rangle$

lemma *iT-Minus-iWeakUntil-conv*: $(P t1. t1 \mathcal{W} t2 (k \ominus I). Q t2) = (P (k - t1). t1 \mathcal{B} t2 (I \downarrow \leq k). Q (k - t2))$

$\langle proof \rangle$

lemma *iT-Div-iWeakUntil-conv*: $(P t1. t1 \mathcal{W} t2 (I \oslash k). Q t2) = (P (t1 \text{ div } k). t1 \mathcal{W} t2 I. Q (t2 \text{ div } k))$

$\langle proof \rangle$

lemma *iT-Plus-iWeakSince-conv*: $(P t1. t1 \mathcal{B} t2 (I \oplus k). Q t2) = (P (t1 + k). t1 \mathcal{B} t2 I. Q (t2 + k))$

$\langle proof \rangle$

lemma *iT-Mult-iWeakSince-conv*: $0 < k \implies (P t1. t1 \mathcal{B} t2 (I \otimes k). Q t2) = (P (t1 * k). t1 \mathcal{B} t2 I. Q (t2 * k))$

$\langle proof \rangle$

lemma *iT-Plus-neg-iWeakSince-conv*: $(P t1. t1 \mathcal{B} t2 (I \oplus - k). Q t2) = (P (t1 - k). t1 \mathcal{B} t2 (I \downarrow \geq k). Q (t2 - k))$

$\langle proof \rangle$

lemma *iT-Minus-iWeakSince-conv*:

$(P t1. t1 \mathcal{B} t2 (k \ominus I). Q t2) = (P (k - t1). t1 \mathcal{W} t2 (I \downarrow \leq k). Q (k - t2))$

thm *iT-Minus-iWeakUntil-conv[symmetric]*

<proof>

lemma *iT-Div-iWeakSince-conv*:

$0 < k \implies (P\ t1.\ t1\ \mathcal{B}\ t2\ (I \odot k).\ Q\ t2) = (P\ (t1\ \text{div}\ k).\ t1\ \mathcal{B}\ t2\ I.\ Q\ (t2\ \text{div}\ k))$

<proof>

Release and Trigger operators

lemma *iT-Plus-iRelease-conv*: $(P\ t1.\ t1\ \mathcal{R}\ t2\ (I \oplus k).\ Q\ t2) = (P\ (t1 + k).\ t1\ \mathcal{R}\ t2\ I.\ Q\ (t2 + k))$

<proof>

lemma *iT-Mult-iRelease-conv*: $(P\ t1.\ t1\ \mathcal{R}\ t2\ (I \otimes k).\ Q\ t2) = (P\ (t1 * k).\ t1\ \mathcal{R}\ t2\ I.\ Q\ (t2 * k))$

<proof>

lemma *iT-Plus-neg-iRelease-conv*: $(P\ t1.\ t1\ \mathcal{R}\ t2\ (I \oplus - k).\ Q\ t2) = (P\ (t1 - k).\ t1\ \mathcal{R}\ t2\ (I \downarrow \geq k).\ Q\ (t2 - k))$

<proof>

lemma *iT-Minus-iRelease-conv*: $(P\ t1.\ t1\ \mathcal{R}\ t2\ (k \ominus I).\ Q\ t2) = (P\ (k - t1).\ t1\ \mathcal{T}\ t2\ (I \downarrow \leq k).\ Q\ (k - t2))$

<proof>

lemma *iT-Div-iRelease-conv*: $(P\ t1.\ t1\ \mathcal{R}\ t2\ (I \odot k).\ Q\ t2) = (P\ (t1\ \text{div}\ k).\ t1\ \mathcal{R}\ t2\ I.\ Q\ (t2\ \text{div}\ k))$

<proof>

lemma *iT-Plus-iTrigger-conv*: $(P\ t1.\ t1\ \mathcal{T}\ t2\ (I \oplus k).\ Q\ t2) = (P\ (t1 + k).\ t1\ \mathcal{T}\ t2\ I.\ Q\ (t2 + k))$

<proof>

lemma *iT-Mult-iTrigger-conv*: $0 < k \implies (P\ t1.\ t1\ \mathcal{T}\ t2\ (I \otimes k).\ Q\ t2) = (P\ (t1 * k).\ t1\ \mathcal{T}\ t2\ I.\ Q\ (t2 * k))$

<proof>

lemma *iT-Plus-neg-iTrigger-conv*: $(P\ t1.\ t1\ \mathcal{T}\ t2\ (I \oplus - k).\ Q\ t2) = (P\ (t1 - k).\ t1\ \mathcal{T}\ t2\ (I \downarrow \geq k).\ Q\ (t2 - k))$

<proof>

lemma *iT-Minus-iTrigger-conv*:

$(P\ t1.\ t1\ \mathcal{T}\ t2\ (k \ominus I).\ Q\ t2) = (P\ (k - t1).\ t1\ \mathcal{R}\ t2\ (I \downarrow \leq k).\ Q\ (k - t2))$

<proof>

lemma *iT-Div-iTrigger-conv*:

$0 < k \implies (P\ t1.\ t1\ \mathcal{T}\ t2\ (I \odot k).\ Q\ t2) = (P\ (t1\ \text{div}\ k).\ t1\ \mathcal{T}\ t2\ I.\ Q\ (t2\ \text{div}\ k))$

<proof>

end